

The L^AT_EX3 Sources

The L^AT_EX Project*

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Abstract

This is the reference documentation for the `expl3` programming environment. The `expl3` modules set up an experimental naming scheme for L^AT_EX commands, which allow the L^AT_EX programmer to systematically name functions and variables, and specify the argument types of functions.

The T_EX and ϵ -T_EX primitives are all given a new name according to these conventions. However, in the main direct use of the primitives is not required or encouraged: the `expl3` modules define an independent low-level L^AT_EX3 programming language.

At present, the `expl3` modules are designed to be loaded on top of L^AT_EX 2 ϵ . In time, a L^AT_EX3 format will be produced based on this code. This allows the code to be used in L^AT_EX 2 ϵ packages *now* while a stand-alone L^AT_EX3 is developed.

While `expl3` is still experimental, the bundle is now regarded as broadly stable. The syntax conventions and functions provided are now ready for wider use. There may still be changes to some functions, but these will be minor when compared to the scope of `expl3`.

New modules will be added to the distributed version of `expl3` as they reach maturity.

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Part I

Introduction to `expl3` and this document

This document is intended to act as a comprehensive reference manual for the `expl3` language. A general guide to the L^AT_EX3 programming language is found in [expl3.pdf](#).

1 Naming functions and variables

L^AT_EX3 does not use `@` as a “letter” for defining internal macros. Instead, the symbols `_` and `:` are used in internal macro names to provide structure. The name of each *function* is divided into logical units using `_`, while `:` separates the *name* of the function from the *argument specifier* (“arg-spec”). This describes the arguments expected by the function. In most cases, each argument is represented by a single letter. The complete list of arg-spec letters for a function is referred to as the *signature* of the function.

Each function name starts with the *module* to which it belongs. Thus apart from a small number of very basic functions, all `expl3` function names contain at least one underscore to divide the module name from the descriptive name of the function. For example, all functions concerned with comma lists are in module `clist` and begin `\clist_`.

Every function must include an argument specifier. For functions which take no arguments, this will be blank and the function name will end `:`. Most functions take one or more arguments, and use the following argument specifiers:

- N and n** These mean *no manipulation*, of a single token for `N` and of a set of tokens given in braces for `n`. Both pass the argument through exactly as given. Usually, if you use a single token for an `n` argument, all will be well.
- c** This means *csname*, and indicates that the argument will be turned into a *csname* before being used. So `\foo:c {ArgumentOne}` will act in the same way as `\foo:N \ArgumentOne`.
- V and v** These mean *value of variable*. The `V` and `v` specifiers are used to get the content of a variable without needing to worry about the underlying T_EX structure containing the data. A `V` argument will be a single token (similar to `N`), for example `\foo:V \MyVariable`; on the other hand, using `v` a *csname* is constructed first, and then the value is recovered, for example `\foo:v {MyVariable}`.
- o** This means *expansion once*. In general, the `V` and `v` specifiers are favoured over `o` for recovering stored information. However, `o` is useful for correctly processing information with delimited arguments.
- x** The `x` specifier stands for *exhaustive expansion*: every token in the argument is fully expanded until only unexpandable ones remain. The T_EX `\edef` primitive carries out this type of expansion. Functions which feature an `x`-type argument are *not* expandable.
- e** The `e` specifier is in many respects identical to `x`, but with a very different implementation. Functions which feature an `e`-type argument may be expandable. The drawback is that `e` is extremely slow (often more than 200 times slower) in older engines, more precisely in non-LuaT_EX engines older than 2019.

- f** The **f** specifier stands for *full expansion*, and in contrast to **x** stops at the first non-expandable token (reading the argument from left to right) without trying to expand it. If this token is a *space token*, it is gobbled, and thus won't be part of the resulting argument. For example, when setting a token list variable (a macro used for storage), the sequence

```
\tl_set:Nn \l_my_a_tl { A }
\tl_set:Nn \l_my_b_tl { B }
\tl_set:Nf \l_my_a_tl { \l_my_a_tl \l_my_b_tl }
```

will leave `\l_my_a_tl` with the content `A\l_my_b_tl`, as `A` cannot be expanded and so terminates expansion before `\l_my_b_tl` is considered.

- T and F** For logic tests, there are the branch specifiers **T** (*true*) and **F** (*false*). Both specifiers treat the input in the same way as **n** (no change), but make the logic much easier to see.
- p** The letter **p** indicates *TeX parameters*. Normally this will be used for delimited functions as `expl3` provides better methods for creating simple sequential arguments.
- w** Finally, there is the **w** specifier for *weird* arguments. This covers everything else, but mainly applies to delimited values (where the argument must be terminated by some specified string).
- D** The **D** stands for **Do not use**. All of the *TeX* primitives are initially `\let` to a **D** name, and some are then given a second name. These functions have no standardized syntax, they are engine dependent and their name can change without warning, thus their use is *strongly discouraged* in package code: programmers should instead use the interfaces documented in [interface3.pdf](#)¹.

Notice that the argument specifier describes how the argument is processed prior to being passed to the underlying function. For example, `\foo:c` will take its argument, convert it to a control sequence and pass it to `\foo:N`.

Variables are named in a similar manner to functions, but begin with a single letter to define the type of variable:

- c** Constant: global parameters whose value should not be changed.
- g** Parameters whose value should only be set globally.
- l** Parameters whose value should only be set locally.

Each variable name is then build up in a similar way to that of a function, typically starting with the module² name and then a descriptive part. Variables end with a short identifier to show the variable type:

clist Comma separated list.

¹If a primitive offers a functionality not yet in the kernel, programmers and users are encouraged to write to the *LaTeX-L* mailing list (<mailto:LATEX-L@listserv.uni-heidelberg.de>) describing their use-case and intended behaviour, so that a possible interface can be discussed. Temporarily, while an interface is not provided, programmers may use the procedure described in the [l3styleguide.pdf](#).

²The module names are not used in case of generic scratch registers defined in the data type modules, e.g., the `int` module contains some scratch variables called `\l_tmpa_int`, `\l_tmpb_int`, and so on. In such a case adding the module name up front to denote the module and in the back to indicate the type, as in `\l_int_tmpa_int` would be very unreadable.

dim “Rigid” lengths.

fp Floating-point values;

int Integer-valued count register.

muskip “Rubber” lengths for use in mathematics.

seq “Sequence”: a data-type used to implement lists (with access at both ends) and stacks.

skip “Rubber” lengths.

str String variables: contain character data.

tl Token list variables: placeholder for a token list.

Applying V-type or v-type expansion to variables of one of the above types is supported, while it is not supported for the following variable types:

bool Either true or false.

box Box register.

coffin A “box with handles” — a higher-level data type for carrying out **box** alignment operations.

flag Integer that can be incremented expandably.

farray Fixed-size array of floating point values.

intarray Fixed-size array of integers.

ior/iow An input or output stream, for reading from or writing to, respectively.

prop Property list: analogue of dictionary or associative arrays in other languages.

regex Regular expression.

1.1 Terminological inexactitude

A word of warning. In this document, and others referring to the `expl3` programming modules, we often refer to “variables” and “functions” as if they were actual constructs from a real programming language. In truth, `TeX` is a macro processor, and functions are simply macros that may or may not take arguments and expand to their replacement text. Many of the common variables are *also* macros, and if placed into the input stream will simply expand to their definition as well — a “function” with no arguments and a “token list variable” are almost the same.³ On the other hand, some “variables” are actually registers that must be initialised and their values set and retrieved with specific functions.

The conventions of the `expl3` code are designed to clearly separate the ideas of “macros that contain data” and “macros that contain code”, and a consistent wrapper is applied to all forms of “data” whether they be macros or actually registers. This means that sometimes we will use phrases like “the function returns a value”, when actually we

³`TeX`nically, functions with no arguments are `\long` while token list variables are not.

just mean “the macro expands to something”. Similarly, the term “execute” might be used in place of “expand” or it might refer to the more specific case of “processing in T_EX’s stomach” (if you are familiar with the T_EXbook parlance).

If in doubt, please ask; chances are we’ve been hasty in writing certain definitions and need to be told to tighten up our terminology.

2 Documentation conventions

This document is typeset with the experimental l3doc class; several conventions are used to help describe the features of the code. A number of conventions are used here to make the documentation clearer.

Each group of related functions is given in a box. For a function with a “user” name, this might read:

```
\ExplSyntaxOn
\ExplSyntaxOff
```

```
\ExplSyntaxOn ... \ExplSyntaxOff
```

The textual description of how the function works would appear here. The syntax of the function is shown in mono-spaced text to the right of the box. In this example, the function takes no arguments and so the name of the function is simply reprinted.

For programming functions, which use `_` and `:` in their name there are a few additional conventions: If two related functions are given with identical names but different argument specifiers, these are termed *variants* of each other, and the latter functions are printed in grey to show this more clearly. They will carry out the same function but will take different types of argument:

```
\seq_new:N
\seq_new:c
```

```
\seq_new:N <sequence>
```

When a number of variants are described, the arguments are usually illustrated only for the base function. Here, `<sequence>` indicates that `\seq_new:N` expects the name of a sequence. From the argument specifier, `\seq_new:c` also expects a sequence name, but as a name rather than as a control sequence. Each argument given in the illustration should be described in the following text.

Fully expandable functions Some functions are fully expandable, which allows them to be used within an `x`-type or `e`-type argument (in plain T_EX terms, inside an `\edef` or `\expanded`), as well as within an `f`-type argument. These fully expandable functions are indicated in the documentation by a star:

```
\cs_to_str:N ☆
```

```
\cs_to_str:N <cs>
```

As with other functions, some text should follow which explains how the function works. Usually, only the star will indicate that the function is expandable. In this case, the function expects a `<cs>`, shorthand for a `<control sequence>`.

Restricted expandable functions A few functions are fully expandable but cannot be fully expanded within an `f`-type argument. In this case a hollow star is used to indicate this:

```
\seq_map_function:NN ☆
```

```
\seq_map_function:NN <seq> <function>
```

Conditional functions Conditional (`if`) functions are normally defined in three variants, with `T`, `F` and `TF` argument specifiers. This allows them to be used for different “true”/“false” branches, depending on which outcome the conditional is being used to test. To indicate this without repetition, this information is given in a shortened form:

<code>\sys_if_engine_xetex:<i><u>TF</u></i> *</code>	<code>\sys_if_engine_xetex:TF {\langle true code \rangle} {\langle false code \rangle}</code>
--	---

The underlining and italic of `TF` indicates that three functions are available:

- `\sys_if_engine_xetex:T`
- `\sys_if_engine_xetex:F`
- `\sys_if_engine_xetex:TF`

Usually, the illustration will use the `TF` variant, and so both $\langle true code \rangle$ and $\langle false code \rangle$ will be shown. The two variant forms `T` and `F` take only $\langle true code \rangle$ and $\langle false code \rangle$, respectively. Here, the star also shows that this function is expandable. With some minor exceptions, *all* conditional functions in the `expl3` modules should be defined in this way.

Variables, constants and so on are described in a similar manner:

`\l_tmpa_tl`

A short piece of text will describe the variable: there is no syntax illustration in this case.

In some cases, the function is similar to one in $\text{\LaTeX} 2_{\epsilon}$ or plain \TeX . In these cases, the text will include an extra “ **\TeX hackers note**” section:

<code>\token_to_str:N *</code>	<code>\token_to_str:N \langle token \rangle</code>
--------------------------------	--

The normal description text.

\TeX hackers note: Detail for the experienced \TeX or $\text{\LaTeX} 2_{\epsilon}$ programmer. In this case, it would point out that this function is the \TeX primitive `\string`.

Changes to behaviour When new functions are added to `expl3`, the date of first inclusion is given in the documentation. Where the documented behaviour of a function changes after it is first introduced, the date of the update will also be given. This means that the programmer can be sure that any release of `expl3` after the date given will contain the function of interest with expected behaviour as described. Note that changes to code internals, including bug fixes, are not recorded in this way *unless* they impact on the expected behaviour.

3 Formal language conventions which apply generally

As this is a formal reference guide for $\text{\LaTeX} 3$ programming, the descriptions of functions are intended to be reasonably “complete”. However, there is also a need to avoid repetition. Formal ideas which apply to general classes of function are therefore summarised here.

For tests which have a `TF` argument specification, the test is evaluated to give a logically `TRUE` or `FALSE` result. Depending on this result, either the $\langle true code \rangle$ or the $\langle false code \rangle$ will be left in the input stream. In the case where the test is expandable,

and a predicate (`_p`) variant is available, the logical value determined by the test is left in the input stream: this will typically be part of a larger logical construct.

4 \TeX concepts not supported by $\text{\LaTeX}3$

The \TeX concept of an “`\outer`” macro is *not supported* at all by $\text{\LaTeX}3$. As such, the functions provided here may break when used on top of $\text{\LaTeX}2_\epsilon$ if `\outer` tokens are used in the arguments.

Part II

The l3bootstrap package

Bootstrap code

1 Using the L^AT_EX3 modules

The modules documented in `source3` are designed to be used on top of L^AT_EX 2_ε and are loaded all as one with the usual `\usepackage{expl3}` or `\RequirePackage{expl3}` instructions.

As the modules use a coding syntax different from standard L^AT_EX 2_ε it provides a few functions for setting it up.

`\ExplSyntaxOn`
`\ExplSyntaxOff`
 Updated: 2011-08-13

`\ExplSyntaxOn` *<code>* `\ExplSyntaxOff`

The `\ExplSyntaxOn` function switches to a category code régime in which spaces are ignored and in which the colon (:) and underscore (_) are treated as “letters”, thus allowing access to the names of code functions and variables. Within this environment, ~ is used to input a space. The `\ExplSyntaxOff` reverts to the document category code régime.

`\ProvidesExplPackage`
`\ProvidesExplClass`
`\ProvidesExplFile`
 Updated: 2017-03-19

`\RequirePackage{expl3}`
`\ProvidesExplPackage` {*<package>*} {*<date>*} {*<version>*} {*<description>*}

These functions act broadly in the same way as the corresponding L^AT_EX 2_ε kernel functions `\ProvidesPackage`, `\ProvidesClass` and `\ProvidesFile`. However, they also implicitly switch `\ExplSyntaxOn` for the remainder of the code with the file. At the end of the file, `\ExplSyntaxOff` will be called to reverse this. (This is the same concept as L^AT_EX 2_ε provides in turning on `\makeatletter` within package and class code.) The *<date>* should be given in the format *<year>/<month>/<day>*. If the *<version>* is given then it will be prefixed with v in the package identifier line.

`\GetIdInfo`
 Updated: 2012-06-04

`\RequirePackage{l3bootstrap}`
`\GetIdInfo` \$Id: *<SVN info field>* \$ {*<description>*}

Extracts all information from a SVN field. Spaces are not ignored in these fields. The information pieces are stored in separate control sequences with `\ExplFileName` for the part of the file name leading up to the period, `\ExplFileDate` for date, `\ExplFileVersion` for version and `\ExplFileDescription` for the description.

To summarize: Every single package using this syntax should identify itself using one of the above methods. Special care is taken so that every package or class file loaded with `\RequirePackage` or similar are loaded with usual L^AT_EX 2_ε category codes and the L^AT_EX3 category code scheme is reloaded when needed afterwards. See implementation for details. If you use the `\GetIdInfo` command you can use the information when loading a package with

```
\ProvidesExplPackage{\ExplFileName}
  {\ExplFileDate}{\ExplFileVersion}{\ExplFileDescription}
```


Part III

The l3names package Namespace for primitives

1 Setting up the L^AT_EX3 programming language

This module is at the core of the L^AT_EX3 programming language. It performs the following tasks:

- defines new names for all T_EX primitives;
- emulate required primitives not provided by default in LuaT_EX;
- switches to the category code régime for programming;

This module is entirely dedicated to primitives (and emulations of these), which should not be used directly within L^AT_EX3 code (outside of “kernel-level” code). As such, the primitives are not documented here: *The T_EXbook*, *T_EX by Topic* and the manuals for pdfT_EX, X_YT_EX, LuaT_EX, pT_EX and upT_EX should be consulted for details of the primitives. These are named `\tex_⟨name⟩:D`, typically based on the primitive’s *⟨name⟩* in pdfT_EX and omitting a leading pdf when the primitive is not related to pdf output.

Part IV

The l3basics package

Basic definitions

As the name suggest this package holds some basic definitions which are needed by most or all other packages in this set.

Here we describe those functions that are used all over the place. With that we mean functions dealing with the construction and testing of control sequences. Furthermore the basic parts of conditional processing are covered; conditional processing dealing with specific data types is described in the modules specific for the respective data types.

1 No operation functions

`\prg_do_nothing: *`

`\prg_do_nothing:`

An expandable function which does nothing at all: leaves nothing in the input stream after a single expansion.

`\scan_stop:`

`\scan_stop:`

A non-expandable function which does nothing. Does not vanish on expansion but produces no typeset output.

2 Grouping material

`\group_begin:`

`\group_begin:`**`\group_end:`**

`\group_end:`

These functions begin and end a group for definition purposes. Assignments are local to groups unless carried out in a global manner. (A small number of exceptions to this rule will be noted as necessary elsewhere in this document.) Each `\group_begin:` must be matched by a `\group_end:`, although this does not have to occur within the same function. Indeed, it is often necessary to start a group within one function and finish it within another, for example when seeking to use non-standard category codes.

`\group_insert_after:N`

`\group_insert_after:N` $\langle token \rangle$

Adds $\langle token \rangle$ to the list of $\langle tokens \rangle$ to be inserted when the current group level ends. The list of $\langle tokens \rangle$ to be inserted is empty at the beginning of a group: multiple applications of `\group_insert_after:N` may be used to build the inserted list one $\langle token \rangle$ at a time. The current group level may be closed by a `\group_end:` function or by a token with category code 2 (close-group), namely a `}` if standard category codes apply.

3 Control sequences and functions

As \TeX is a macro language, creating new functions means creating macros. At point of use, a function is replaced by the replacement text (“code”) in which each parameter in the code ($\#1$, $\#2$, *etc.*) is replaced the appropriate arguments absorbed by the function. In the following, *code* is therefore used as a shorthand for “replacement text”.

Functions which are not “protected” are fully expanded inside an \mathbf{x} expansion. In contrast, “protected” functions are not expanded within \mathbf{x} expansions.

3.1 Defining functions

Functions can be created with no requirement that they are declared first (in contrast to variables, which must always be declared). Declaring a function before setting up the code means that the name chosen is checked and an error raised if it is already in use. The name of a function can be checked at the point of definition using the `\cs_new...` functions: this is recommended for all functions which are defined for the first time.

There are three ways to define new functions. All classes define a function to expand to the substitution text. Within the substitution text the actual parameters are substituted for the formal parameters ($\#1$, $\#2$, ...).

new Create a new function with the **new** scope, such as `\cs_new:Npn`. The definition is global and results in an error if it is already defined.

set Create a new function with the **set** scope, such as `\cs_set:Npn`. The definition is restricted to the current \TeX group and does not result in an error if the function is already defined.

gset Create a new function with the **gset** scope, such as `\cs_gset:Npn`. The definition is global and does not result in an error if the function is already defined.

Within each set of scope there are different ways to define a function. The differences depend on restrictions on the actual parameters and the expandability of the resulting function.

nopar Create a new function with the **nopar** restriction, such as `\cs_set_nopar:Npn`. The parameter may not contain `\par` tokens.

protected Create a new function with the **protected** restriction, such as `\cs_set_protected:Npn`. The parameter may contain `\par` tokens but the function will not expand within an \mathbf{x} -type or \mathbf{e} -type expansion.

Finally, the functions in Subsections 3.2 and 3.3 are primarily meant to define *base functions* only. Base functions can only have the following argument specifiers:

N and n No manipulation.

T and F Functionally equivalent to **n** (you are actually encouraged to use the family of `\prg_new_conditional:` functions described in Section 1).

p and w These are special cases.

The `\cs_new:` functions below (and friends) do not stop you from using other argument specifiers in your function names, but they do not handle expansion for you. You should define the base function and then use `\cs_generate_variant:Nn` to generate custom variants as described in Section 2.

3.2 Defining new functions using parameter text

<code>\cs_new:Npn</code>	<code>\cs_new:Npn <function> <parameters> {<code>}</code>
<code>\cs_new:cpn</code>	Creates <code><function></code> to expand to <code><code></code> as replacement text. Within the <code><code></code> , the
<code>\cs_new:Npx</code>	<code><parameters></code> (#1, #2, etc.) will be replaced by those absorbed by the function. The
<code>\cs_new:cpx</code>	definition is global and an error results if the <code><function></code> is already defined.

<code>\cs_new_nopar:Npn</code>	<code>\cs_new_nopar:Npn <function> <parameters> {<code>}</code>
<code>\cs_new_nopar:cpn</code>	Creates <code><function></code> to expand to <code><code></code> as replacement text. Within the <code><code></code> , the
<code>\cs_new_nopar:Npx</code>	<code><parameters></code> (#1, #2, etc.) will be replaced by those absorbed by the function. When the
<code>\cs_new_nopar:cpx</code>	<code><function></code> is used the <code><parameters></code> absorbed cannot contain <code>\par</code> tokens. The definition
	is global and an error results if the <code><function></code> is already defined.

<code>\cs_new_protected:Npn</code>	<code>\cs_new_protected:Npn <function> <parameters> {<code>}</code>
<code>\cs_new_protected:cpn</code>	Creates <code><function></code> to expand to <code><code></code> as replacement text. Within the <code><code></code> , the
<code>\cs_new_protected:Npx</code>	<code><parameters></code> (#1, #2, etc.) will be replaced by those absorbed by the function. The
<code>\cs_new_protected:cpx</code>	<code><function></code> will not expand within an x-type argument. The definition is global and an
	error results if the <code><function></code> is already defined.

<code>\cs_new_protected_nopar:Npn</code>	<code>\cs_new_protected_nopar:Npn <function> <parameters> {<code>}</code>
<code>\cs_new_protected_nopar:cpn</code>	
<code>\cs_new_protected_nopar:Npx</code>	
<code>\cs_new_protected_nopar:cpx</code>	

Creates `<function>` to expand to `<code>` as replacement text. Within the `<code>`, the `<parameters>` (#1, #2, etc.) will be replaced by those absorbed by the function. When the `<function>` is used the `<parameters>` absorbed cannot contain `\par` tokens. The `<function>` will not expand within an x-type or e-type argument. The definition is global and an error results if the `<function>` is already defined.

<code>\cs_set:Npn</code>	<code>\cs_set:Npn <function> <parameters> {<code>}</code>
<code>\cs_set:cpn</code>	Sets <code><function></code> to expand to <code><code></code> as replacement text. Within the <code><code></code> , the
<code>\cs_set:Npx</code>	<code><parameters></code> (#1, #2, etc.) will be replaced by those absorbed by the function. The
<code>\cs_set:cpx</code>	assignment of a meaning to the <code><function></code> is restricted to the current TeX group level.

<code>\cs_set_nopar:Npn</code>	<code>\cs_set_nopar:Npn <function> <parameters> {<code>}</code>
<code>\cs_set_nopar:cpn</code>	Sets <code><function></code> to expand to <code><code></code> as replacement text. Within the <code><code></code> , the
<code>\cs_set_nopar:Npx</code>	<code><parameters></code> (#1, #2, etc.) will be replaced by those absorbed by the function. When the
<code>\cs_set_nopar:cpx</code>	<code><function></code> is used the <code><parameters></code> absorbed cannot contain <code>\par</code> tokens. The assignment
	of a meaning to the <code><function></code> is restricted to the current TeX group level.

<code>\cs_set_protected:Npn</code>	<code>\cs_set_protected:Npn <function> <parameters> {<code>}</code>
<code>\cs_set_protected:cpn</code>	Sets <code><function></code> to expand to <code><code></code> as replacement text. Within the <code><code></code> , the
<code>\cs_set_protected:Npx</code>	<code><parameters></code> (#1, #2, etc.) will be replaced by those absorbed by the function. The
<code>\cs_set_protected:cpx</code>	assignment of a meaning to the <code><function></code> is restricted to the current TeX group level.
	The <code><function></code> will not expand within an x-type or e-type argument.

<code>\cs_set_protected_nopar:Npn</code>	<code>\cs_set_protected_nopar:Npn <function> <parameters> {<code>}</code>
<code>\cs_set_protected_nopar:cpn</code>	
<code>\cs_set_protected_nopar:Npx</code>	
<code>\cs_set_protected_nopar:cpx</code>	

Sets $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the $\langle parameters \rangle$ ($\#1$, $\#2$, *etc.*) will be replaced by those absorbed by the function. When the $\langle function \rangle$ is used the $\langle parameters \rangle$ absorbed cannot contain `\par` tokens. The assignment of a meaning to the $\langle function \rangle$ is restricted to the current \TeX group level. The $\langle function \rangle$ will not expand within an **x**-type or **e**-type argument.

<code>\cs_gset:Npn</code>	<code>\cs_gset:Npn <function> <parameters> {<code>}</code>
<code>\cs_gset:cpn</code>	
<code>\cs_gset:Npx</code>	
<code>\cs_gset:cpx</code>	

Globally sets $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the $\langle parameters \rangle$ ($\#1$, $\#2$, *etc.*) will be replaced by those absorbed by the function. The assignment of a meaning to the $\langle function \rangle$ is *not* restricted to the current \TeX group level: the assignment is global.

<code>\cs_gset_nopar:Npn</code>	<code>\cs_gset_nopar:Npn <function> <parameters> {<code>}</code>
<code>\cs_gset_nopar:cpn</code>	
<code>\cs_gset_nopar:Npx</code>	
<code>\cs_gset_nopar:cpx</code>	

Globally sets $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the $\langle parameters \rangle$ ($\#1$, $\#2$, *etc.*) will be replaced by those absorbed by the function. When the $\langle function \rangle$ is used the $\langle parameters \rangle$ absorbed cannot contain `\par` tokens. The assignment of a meaning to the $\langle function \rangle$ is *not* restricted to the current \TeX group level: the assignment is global.

<code>\cs_gset_protected:Npn</code>	<code>\cs_gset_protected:Npn <function> <parameters> {<code>}</code>
<code>\cs_gset_protected:cpn</code>	
<code>\cs_gset_protected:Npx</code>	
<code>\cs_gset_protected:cpx</code>	

Globally sets $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the $\langle parameters \rangle$ ($\#1$, $\#2$, *etc.*) will be replaced by those absorbed by the function. The assignment of a meaning to the $\langle function \rangle$ is *not* restricted to the current \TeX group level: the assignment is global. The $\langle function \rangle$ will not expand within an **x**-type or **e**-type argument.

<code>\cs_gset_protected_nopar:Npn</code>	<code>\cs_gset_protected_nopar:Npn <function> <parameters> {<code>}</code>
<code>\cs_gset_protected_nopar:cpn</code>	
<code>\cs_gset_protected_nopar:Npx</code>	
<code>\cs_gset_protected_nopar:cpx</code>	

Globally sets $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the $\langle parameters \rangle$ ($\#1$, $\#2$, *etc.*) will be replaced by those absorbed by the function. When the $\langle function \rangle$ is used the $\langle parameters \rangle$ absorbed cannot contain `\par` tokens. The assignment of a meaning to the $\langle function \rangle$ is *not* restricted to the current \TeX group level: the assignment is global. The $\langle function \rangle$ will not expand within an **x**-type argument.

3.3 Defining new functions using the signature

<code>\cs_new:Nn</code>	<code>\cs_new:Nn <function> {<code>}</code>
<code>\cs_new:(cn Nx cx)</code>	

Creates $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the number of $\langle parameters \rangle$ is detected automatically from the function signature. These $\langle parameters \rangle$ ($\#1$, $\#2$, *etc.*) will be replaced by those absorbed by the function. The definition is global and an error results if the $\langle function \rangle$ is already defined.

<hr/> <code>\cs_new_nopar:Nn</code> <code>\cs_new_nopar:(cn Nx cx)</code> <hr/>	<code>\cs_new_nopar:Nn <function> {<code>}</code> <p>Creates $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the number of $\langle parameters \rangle$ is detected automatically from the function signature. These $\langle parameters \rangle$ ($\#1$, $\#2$, <i>etc.</i>) will be replaced by those absorbed by the function. When the $\langle function \rangle$ is used the $\langle parameters \rangle$ absorbed cannot contain <code>\par</code> tokens. The definition is global and an error results if the $\langle function \rangle$ is already defined.</p>
<hr/> <code>\cs_new_protected:Nn</code> <code>\cs_new_protected:(cn Nx cx)</code> <hr/>	<code>\cs_new_protected:Nn <function> {<code>}</code> <p>Creates $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the number of $\langle parameters \rangle$ is detected automatically from the function signature. These $\langle parameters \rangle$ ($\#1$, $\#2$, <i>etc.</i>) will be replaced by those absorbed by the function. The $\langle function \rangle$ will not expand within an x-type argument. The definition is global and an error results if the $\langle function \rangle$ is already defined.</p>
<hr/> <code>\cs_new_protected_nopar:Nn</code> <code>\cs_new_protected_nopar:(cn Nx cx)</code> <hr/>	<code>\cs_new_protected_nopar:Nn <function> {<code>}</code> <p>Creates $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the number of $\langle parameters \rangle$ is detected automatically from the function signature. These $\langle parameters \rangle$ ($\#1$, $\#2$, <i>etc.</i>) will be replaced by those absorbed by the function. When the $\langle function \rangle$ is used the $\langle parameters \rangle$ absorbed cannot contain <code>\par</code> tokens. The $\langle function \rangle$ will not expand within an x-type or e-type argument. The definition is global and an error results if the $\langle function \rangle$ is already defined.</p>
<hr/> <code>\cs_set:Nn</code> <code>\cs_set:(cn Nx cx)</code> <hr/>	<code>\cs_set:Nn <function> {<code>}</code> <p>Sets $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the number of $\langle parameters \rangle$ is detected automatically from the function signature. These $\langle parameters \rangle$ ($\#1$, $\#2$, <i>etc.</i>) will be replaced by those absorbed by the function. The assignment of a meaning to the $\langle function \rangle$ is restricted to the current \TeX group level.</p>
<hr/> <code>\cs_set_nopar:Nn</code> <code>\cs_set_nopar:(cn Nx cx)</code> <hr/>	<code>\cs_set_nopar:Nn <function> {<code>}</code> <p>Sets $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the number of $\langle parameters \rangle$ is detected automatically from the function signature. These $\langle parameters \rangle$ ($\#1$, $\#2$, <i>etc.</i>) will be replaced by those absorbed by the function. When the $\langle function \rangle$ is used the $\langle parameters \rangle$ absorbed cannot contain <code>\par</code> tokens. The assignment of a meaning to the $\langle function \rangle$ is restricted to the current \TeX group level.</p>
<hr/> <code>\cs_set_protected:Nn</code> <code>\cs_set_protected:(cn Nx cx)</code> <hr/>	<code>\cs_set_protected:Nn <function> {<code>}</code> <p>Sets $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the number of $\langle parameters \rangle$ is detected automatically from the function signature. These $\langle parameters \rangle$ ($\#1$, $\#2$, <i>etc.</i>) will be replaced by those absorbed by the function. The $\langle function \rangle$ will not expand within an x-type argument. The assignment of a meaning to the $\langle function \rangle$ is restricted to the current \TeX group level.</p>

<code>\cs_set_protected_nopar:Nn</code>	<code>\cs_set_protected_nopar:Nn <function> {<code>}</code>
<code>\cs_set_protected_nopar:(cn Nx cx)</code>	

Sets $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the number of $\langle parameters \rangle$ is detected automatically from the function signature. These $\langle parameters \rangle$ ($\#1$, $\#2$, *etc.*) will be replaced by those absorbed by the function. When the $\langle function \rangle$ is used the $\langle parameters \rangle$ absorbed cannot contain `\par` tokens. The $\langle function \rangle$ will not expand within an *x*-type or *e*-type argument. The assignment of a meaning to the $\langle function \rangle$ is restricted to the current T_EX group level.

<code>\cs_gset:Nn</code>	<code>\cs_gset:Nn <function> {<code>}</code>
<code>\cs_gset:(cn Nx cx)</code>	

Sets $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the number of $\langle parameters \rangle$ is detected automatically from the function signature. These $\langle parameters \rangle$ ($\#1$, $\#2$, *etc.*) will be replaced by those absorbed by the function. The assignment of a meaning to the $\langle function \rangle$ is global.

<code>\cs_gset_nopar:Nn</code>	<code>\cs_gset_nopar:Nn <function> {<code>}</code>
<code>\cs_gset_nopar:(cn Nx cx)</code>	

Sets $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the number of $\langle parameters \rangle$ is detected automatically from the function signature. These $\langle parameters \rangle$ ($\#1$, $\#2$, *etc.*) will be replaced by those absorbed by the function. When the $\langle function \rangle$ is used the $\langle parameters \rangle$ absorbed cannot contain `\par` tokens. The assignment of a meaning to the $\langle function \rangle$ is global.

<code>\cs_gset_protected:Nn</code>	<code>\cs_gset_protected:Nn <function> {<code>}</code>
<code>\cs_gset_protected:(cn Nx cx)</code>	

Sets $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the number of $\langle parameters \rangle$ is detected automatically from the function signature. These $\langle parameters \rangle$ ($\#1$, $\#2$, *etc.*) will be replaced by those absorbed by the function. The $\langle function \rangle$ will not expand within an *x*-type argument. The assignment of a meaning to the $\langle function \rangle$ is global.

<code>\cs_gset_protected_nopar:Nn</code>	<code>\cs_gset_protected_nopar:Nn <function> {<code>}</code>
<code>\cs_gset_protected_nopar:(cn Nx cx)</code>	

Sets $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the number of $\langle parameters \rangle$ is detected automatically from the function signature. These $\langle parameters \rangle$ ($\#1$, $\#2$, *etc.*) will be replaced by those absorbed by the function. When the $\langle function \rangle$ is used the $\langle parameters \rangle$ absorbed cannot contain `\par` tokens. The $\langle function \rangle$ will not expand within an *x*-type or *e*-type argument. The assignment of a meaning to the $\langle function \rangle$ is global.

<code>\cs_generate_from_arg_count:NNnn</code>	<code>\cs_generate_from_arg_count:NNnn <function> <creator> {<number>}</code>
<code>\cs_generate_from_arg_count:(cNnn Ncnn)</code>	<code>{<code>}</code>

Updated: 2012-01-14

Uses the $\langle creator \rangle$ function (which should have signature Npn , for example `\cs_new:Npn`) to define a $\langle function \rangle$ which takes $\langle number \rangle$ arguments and has $\langle code \rangle$ as replacement text. The $\langle number \rangle$ of arguments is an integer expression, evaluated as detailed for `\int_eval:n`.

3.4 Copying control sequences

Control sequences (not just functions as defined above) can be set to have the same meaning using the functions described here. Making two control sequences equivalent means that the second control sequence is a *copy* of the first (rather than a pointer to it). Thus the old and new control sequence are not tied together: changes to one are not reflected in the other.

In the following text “cs” is used as an abbreviation for “control sequence”.

```
\cs_new_eq:NN
\cs_new_eq:(Nc|cN|cc)
```

```
\cs_new_eq:NN <cs1> <cs2>
\cs_new_eq:NN <cs1> <token>
```

Globally creates $\langle control\ sequence_1 \rangle$ and sets it to have the same meaning as $\langle control\ sequence_2 \rangle$ or $\langle token \rangle$. The second control sequence may subsequently be altered without affecting the copy.

```
\cs_set_eq:NN
\cs_set_eq:(Nc|cN|cc)
```

```
\cs_set_eq:NN <cs1> <cs2>
\cs_set_eq:NN <cs1> <token>
```

Sets $\langle control\ sequence_1 \rangle$ to have the same meaning as $\langle control\ sequence_2 \rangle$ (or $\langle token \rangle$). The second control sequence may subsequently be altered without affecting the copy. The assignment of a meaning to the $\langle control\ sequence_1 \rangle$ is restricted to the current \TeX group level.

```
\cs_gset_eq:NN
\cs_gset_eq:(Nc|cN|cc)
```

```
\cs_gset_eq:NN <cs1> <cs2>
\cs_gset_eq:NN <cs1> <token>
```

Globally sets $\langle control\ sequence_1 \rangle$ to have the same meaning as $\langle control\ sequence_2 \rangle$ (or $\langle token \rangle$). The second control sequence may subsequently be altered without affecting the copy. The assignment of a meaning to the $\langle control\ sequence_1 \rangle$ is *not* restricted to the current \TeX group level: the assignment is global.

3.5 Deleting control sequences

There are occasions where control sequences need to be deleted. This is handled in a very simple manner.

```
\cs_undefine:N
\cs_undefine:c
```

```
\cs_undefine:N <control\ sequence>
```

Sets $\langle control\ sequence \rangle$ to be globally undefined.

Updated: 2011-09-15

3.6 Showing control sequences

```
\cs_meaning:N ★
\cs_meaning:c ★
```

```
\cs_meaning:N <control\ sequence>
```

This function expands to the *meaning* of the $\langle control\ sequence \rangle$ control sequence. For a macro, this includes the $\langle replacement\ text \rangle$.

Updated: 2011-12-22

\TeX hackers note: This is \TeX ’s $\backslash meaning$ primitive. For tokens that are not control sequences, it is more logical to use $\backslash token_to_meaning:N$. The *c* variant correctly reports undefined arguments.

`\cs_show:N`
`\cs_show:c`

Updated: 2017-02-14

`\cs_show:N` $\langle control\ sequence \rangle$
Displays the definition of the $\langle control\ sequence \rangle$ on the terminal.

T_EXhackers note: This is similar to the T_EX primitive `\show`, wrapped to a fixed number of characters per line.

`\cs_log:N`
`\cs_log:c`

New: 2014-08-22
Updated: 2017-02-14

`\cs_log:N` $\langle control\ sequence \rangle$
Writes the definition of the $\langle control\ sequence \rangle$ in the log file. See also `\cs_show:N` which displays the result in the terminal.

3.7 Converting to and from control sequences

`\use:c` ★

`\use:c` $\{ \langle control\ sequence\ name \rangle \}$

Expands the $\langle control\ sequence\ name \rangle$ until only characters remain, and then converts this into a control sequence. This process requires two expansions. As in other `c`-type arguments the $\langle control\ sequence\ name \rangle$ must, when fully expanded, consist of character tokens, typically a mixture of category code 10 (space), 11 (letter) and 12 (other).

T_EXhackers note: Protected macros that appear in a `c`-type argument are expanded despite being protected; `\exp_not:n` also has no effect. An internal error occurs if non-characters or active characters remain after full expansion, as the conversion to a control sequence is not possible.

As an example of the `\use:c` function, both

`\use:c { a b c }`

and

```
\tl_new:N \l_my_tl
\tl_set:Nn \l_my_tl { a b c }
\use:c { \tl_use:N \l_my_tl }
```

would be equivalent to

`\abc`

after two expansions of `\use:c`.

`\cs_if_exist_use:N` ★
`\cs_if_exist_use:c` ★
`\cs_if_exist_use:NTF` ★
`\cs_if_exist_use:cTF` ★

New: 2012-11-10

`\cs_if_exist_use:N` $\langle control\ sequence \rangle$
`\cs_if_exist_use:NTF` $\langle control\ sequence \rangle$ $\{ \langle true\ code \rangle \}$ $\{ \langle false\ code \rangle \}$

Tests whether the $\langle control\ sequence \rangle$ is currently defined according to the conditional `\cs_if_exist:NTF` (whether as a function or another control sequence type), and if it is inserts the $\langle control\ sequence \rangle$ into the input stream followed by the $\langle true\ code \rangle$. Otherwise the $\langle false\ code \rangle$ is used.

<code>\cs:w</code>	★	<code>\cs:w</code> \langle <i>control sequence name</i> \rangle <code>\cs_end:</code>
<code>\cs_end:</code>	★	

Converts the given \langle *control sequence name* \rangle into a single control sequence token. This process requires one expansion. The content for \langle *control sequence name* \rangle may be literal material or from other expandable functions. The \langle *control sequence name* \rangle must, when fully expanded, consist of character tokens which are not active: typically of category code 10 (space), 11 (letter) or 12 (other), or a mixture of these.

T_EXhackers note: These are the T_EX primitives `\csname` and `\endcsname`.

As an example of the `\cs:w` and `\cs_end:` functions, both

`\cs:w a b c \cs_end:`

and

```
\tl_new:N \l_my_tl
\tl_set:Nn \l_my_tl { a b c }
\cs:w \tl_use:N \l_my_tl \cs_end:
```

would be equivalent to

`\abc`

after one expansion of `\cs:w`.

<code>\cs_to_str:N</code>	★	<code>\cs_to_str:N</code> \langle <i>control sequence</i> \rangle
---------------------------	---	---

Converts the given \langle *control sequence* \rangle into a series of characters with category code 12 (other), except spaces, of category code 10. The result does *not* include the current escape token, contrarily to `\token_to_str:N`. Full expansion of this function requires exactly 2 expansion steps, and so an x-type or e-type expansion, or two o-type expansions are required to convert the \langle *control sequence* \rangle to a sequence of characters in the input stream. In most cases, an f-expansion is correct as well, but this loses a space at the start of the result.

4 Analysing control sequences

<code>\cs_split_function:N</code>	★	<code>\cs_split_function:N</code> \langle <i>function</i> \rangle
-----------------------------------	---	---

New: 2018-04-06

Splits the \langle *function* \rangle into the \langle *name* \rangle (*i.e.* the part before the colon) and the \langle *signature* \rangle (*i.e.* after the colon). This information is then placed in the input stream in three parts: the \langle *name* \rangle , the \langle *signature* \rangle and a logic token indicating if a colon was found (to differentiate variables from function names). The \langle *name* \rangle does not include the escape character, and both the \langle *name* \rangle and \langle *signature* \rangle are made up of tokens with category code 12 (other).

The next three functions decompose T_EX macros into their constituent parts: if the \langle *token* \rangle passed is not a macro then no decomposition can occur. In the latter case, all three functions leave `\scan_stop:` in the input stream.

\cs_prefix_spec:N ★

New: 2019-02-27

\cs_prefix_spec:N $\langle token \rangle$

If the $\langle token \rangle$ is a macro, this function leaves the applicable TeX prefixes in input stream as a string of tokens of category code 12 (with spaces having category code 10). Thus for example

```
\cs_set:Npn \next:nn #1#2 { x #1~y #2 }
\cs_prefix_spec:N \next:nn
```

leaves `\long` in the input stream. If the $\langle token \rangle$ is not a macro then `\scan_stop:` is left in the input stream.

TeXhackers note: The prefix can be empty, `\long`, `\protected` or `\protected\long` with backslash replaced by the current escape character.

\cs_argument_spec:N ★

New: 2019-02-27

\cs_argument_spec:N $\langle token \rangle$

If the $\langle token \rangle$ is a macro, this function leaves the primitive TeX argument specification in input stream as a string of character tokens of category code 12 (with spaces having category code 10). Thus for example

```
\cs_set:Npn \next:nn #1#2 { x #1 y #2 }
\cs_argument_spec:N \next:nn
```

leaves `#1#2` in the input stream. If the $\langle token \rangle$ is not a macro then `\scan_stop:` is left in the input stream.

TeXhackers note: If the argument specification contains the string `->`, then the function produces incorrect results.

\cs_replacement_spec:N ★

New: 2019-02-27

\cs_replacement_spec:N $\langle token \rangle$

If the $\langle token \rangle$ is a macro, this function leaves the replacement text in input stream as a string of character tokens of category code 12 (with spaces having category code 10). Thus for example

```
\cs_set:Npn \next:nn #1#2 { x #1~y #2 }
\cs_replacement_spec:N \next:nn
```

leaves `x#1~y#2` in the input stream. If the $\langle token \rangle$ is not a macro then `\scan_stop:` is left in the input stream.

TeXhackers note: If the argument specification contains the string `->`, then the function produces incorrect results.

5 Using or removing tokens and arguments

Tokens in the input can be read and used or read and discarded. If one or more tokens are wrapped in braces then when absorbing them the outer set is removed. At the same time, the category code of each token is set when the token is read by a function (if it

is read more than once, the category code is determined by the situation in force when first function absorbs the token).

<code>\use:n</code>	*	<code>\use:n</code>	<code>{\langle group_1 \rangle}</code>
<code>\use:nn</code>	*	<code>\use:nn</code>	<code>{\langle group_1 \rangle} {\langle group_2 \rangle}</code>
<code>\use:nnn</code>	*	<code>\use:nnn</code>	<code>{\langle group_1 \rangle} {\langle group_2 \rangle} {\langle group_3 \rangle}</code>
<code>\use:nnnn</code>	*	<code>\use:nnnn</code>	<code>{\langle group_1 \rangle} {\langle group_2 \rangle} {\langle group_3 \rangle} {\langle group_4 \rangle}</code>

As illustrated, these functions absorb between one and four arguments, as indicated by the argument specifier. The braces surrounding each argument are removed and the remaining tokens are left in the input stream. The category code of these tokens is also fixed by this process (if it has not already been by some other absorption). All of these functions require only a single expansion to operate, so that one expansion of

`\use:nn { abc } { { def } }`

results in the input stream containing

`abc { def }`

i.e. only the outer braces are removed.

T_EXhackers note: The `\use:n` function is equivalent to L^AT_EX 2_ε's `\@firstofone`.

<code>\use_i:nn</code>	*	<code>\use_i:nn</code>	<code>{\langle arg_1 \rangle} {\langle arg_2 \rangle}</code>
<code>\use_ii:nn</code>	*		

These functions absorb two arguments from the input stream. The function `\use_i:nn` discards the second argument, and leaves the content of the first argument in the input stream. `\use_ii:nn` discards the first argument and leaves the content of the second argument in the input stream. The category code of these tokens is also fixed (if it has not already been by some other absorption). A single expansion is needed for the functions to take effect.

T_EXhackers note: These are equivalent to L^AT_EX 2_ε's `\@firstoftwo` and `\@secondoftwo`.

<code>\use_i:nnn</code>	*	<code>\use_i:nnn</code>	<code>{\langle arg_1 \rangle} {\langle arg_2 \rangle} {\langle arg_3 \rangle}</code>
<code>\use_ii:nnn</code>	*		
<code>\use_iii:nnn</code>	*		

These functions absorb three arguments from the input stream. The function `\use_i:nnn` discards the second and third arguments, and leaves the content of the first argument in the input stream. `\use_ii:nnn` and `\use_iii:nnn` work similarly, leaving the content of second or third arguments in the input stream, respectively. The category code of these tokens is also fixed (if it has not already been by some other absorption). A single expansion is needed for the functions to take effect.

<code>\use_i:nnnn</code>	*	<code>\use_i:nnnn</code>	<code>{\langle arg_1 \rangle} {\langle arg_2 \rangle} {\langle arg_3 \rangle} {\langle arg_4 \rangle}</code>
<code>\use_ii:nnnn</code>	*		
<code>\use_iii:nnnn</code>	*		
<code>\use_iv:nnnn</code>	*		

These functions absorb four arguments from the input stream. The function `\use_i:nnnn` discards the second, third and fourth arguments, and leaves the content of the first argument in the input stream. `\use_ii:nnnn`, `\use_iii:nnnn` and `\use_iv:nnnn` work similarly, leaving the content of second, third or fourth arguments in the input stream, respectively. The category code of these tokens is also fixed (if it has not already been by some other absorption). A single expansion is needed for the functions to take effect.

<code>\use_i_ii:nnn</code>	★	<code>\use_i_ii:nnn {\langle arg_1 \rangle} {\langle arg_2 \rangle} {\langle arg_3 \rangle}</code>
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This function absorbs three arguments and leaves the content of the first and second in the input stream. The category code of these tokens is also fixed (if it has not already been by some other absorption). A single expansion is needed for the function to take effect. An example:

```
\use_i_ii:nnn { abc } { { def } } { ghi }
```

results in the input stream containing

```
abc { def }
```

i.e. the outer braces are removed and the third group is removed.

<code>\use_ii_i:nn</code>	★	<code>\use_ii_i:nn {\langle arg_1 \rangle} {\langle arg_2 \rangle}</code>
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New: 2019-06-02

This function absorbs two arguments and leaves the content of the second and first in the input stream. The category code of these tokens is also fixed (if it has not already been by some other absorption). A single expansion is needed for the function to take effect.

<code>\use_none:n</code>	★	<code>\use_none:n {\langle group_1 \rangle}</code>
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<code>\use_none:nn</code>	★
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<code>\use_none:nnn</code>	★
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<code>\use_none:nnnn</code>	★
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<code>\use_none:nnnnn</code>	★
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<code>\use_none:nnnnnn</code>	★
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<code>\use_none:nnnnnnn</code>	★
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<code>\use_none:nnnnnnnn</code>	★
---------------------------------	---

<code>\use_none:nnnnnnnnn</code>	★
----------------------------------	---

These functions absorb between one and nine groups from the input stream, leaving nothing on the resulting input stream. These functions work after a single expansion. One or more of the `n` arguments may be an unbraced single token (*i.e.* an `N` argument).

TeXhackers note: These are equivalent to L^AT_EX 2_ε's `\@gobble`, `\@gobbletwo`, *etc.*

<code>\use:e</code>	★	<code>\use:e {\langle expandable tokens \rangle}</code>
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New: 2018-06-18

Fully expands the `\langle token list \rangle` in an `x`-type manner, *but* the function remains fully expandable, and parameter character (usually `#`) need not be doubled.

TeXhackers note: `\use:e` is a wrapper around the primitive `\expanded` where it is available: it requires two expansions to complete its action. When `\expanded` is not available this function is very slow.

<code>\use:x</code>		<code>\use:x {\langle expandable tokens \rangle}</code>
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Updated: 2011-12-31

Fully expands the `\langle expandable tokens \rangle` and inserts the result into the input stream at the current location. Any hash characters (`#`) in the argument must be doubled.

5.1 Selecting tokens from delimited arguments

A different kind of function for selecting tokens from the token stream are those that use delimited arguments.

<code>\use_none_delimit_by_q_nil:w</code>	<code>*</code>	<code>\use_none_delimit_by_q_nil:w <balanced text> \q_nil</code>
<code>\use_none_delimit_by_q_stop:w</code>	<code>*</code>	<code>\use_none_delimit_by_q_stop:w <balanced text> \q_stop</code>
<code>\use_none_delimit_by_q_recursion_stop:w</code>	<code>*</code>	<code>\use_none_delimit_by_q_recursion_stop:w <balanced text></code>
		<code>\q_recursion_stop</code>

Absorb the *<balanced text>* from the input stream delimited by the marker given in the function name, leaving nothing in the input stream.

<code>\use_i_delimit_by_q_nil:nw</code>	<code>*</code>	<code>\use_i_delimit_by_q_nil:nw {<inserted tokens>} <balanced text></code>
<code>\use_i_delimit_by_q_stop:nw</code>	<code>*</code>	<code>\q_nil</code>
<code>\use_i_delimit_by_q_recursion_stop:nw</code>	<code>*</code>	<code>\use_i_delimit_by_q_stop:nw {<inserted tokens>} <balanced</code>
		<code>text> \q_stop</code>
		<code>\use_i_delimit_by_q_recursion_stop:nw {<inserted tokens>}</code>
		<code><balanced text> \q_recursion_stop</code>

Absorb the *<balanced text>* from the input stream delimited by the marker given in the function name, leaving *<inserted tokens>* in the input stream for further processing.

6 Predicates and conditionals

L^AT_EX3 has three concepts for conditional flow processing:

Branching conditionals Functions that carry out a test and then execute, depending on its result, either the code supplied as the *<true code>* or the *<false code>*. These arguments are denoted with T and F, respectively. An example would be

`\cs_if_free:cTF {abc} {<true code>} {<false code>}`

a function that turns the first argument into a control sequence (since it's marked as c) then checks whether this control sequence is still free and then depending on the result carries out the code in the second argument (true case) or in the third argument (false case).

These type of functions are known as “conditionals”; whenever a TF function is defined it is usually accompanied by T and F functions as well. These are provided for convenience when the branch only needs to go a single way. Package writers are free to choose which types to define but the kernel definitions always provide all three versions.

Important to note is that these branching conditionals with *<true code>* and/or *<false code>* are always defined in a way that the code of the chosen alternative can operate on following tokens in the input stream.

These conditional functions may or may not be fully expandable, but if they are expandable they are accompanied by a “predicate” for the same test as described below.

Predicates “Predicates” are functions that return a special type of boolean value which can be tested by the boolean expression parser. All functions of this type are expandable and have names that end with `_p` in the description part. For example,

`\cs_if_free_p:N`

would be a predicate function for the same type of test as the conditional described above. It would return “true” if its argument (a single token denoted by N) is still free for definition. It would be used in constructions like

```

\bool_if:nTF {
  \cs_if_free_p:N \l_tmpz_tl || \cs_if_free_p:N \g_tmpz_tl
} {\true code} {\false code}

```

For each predicate defined, a “branching conditional” also exists that behaves like a conditional described above.

Primitive conditionals There is a third variety of conditional, which is the original concept used in plain T_EX and L^AT_EX 2_ε. Their use is discouraged in expl3 (although still used in low-level definitions) because they are more fragile and in many cases require more expansion control (hence more code) than the two types of conditionals described above.

```

\c_true_bool
\c_false_bool

```

Constants that represent `true` and `false`, respectively. Used to implement predicates.

6.1 Tests on control sequences

```

\cs_if_eq_p:NN *
\cs_if_eq:NNTF *

```

```

\cs_if_eq_p:NN <cs1> <cs2>
\cs_if_eq:NNTF <cs1> <cs2> {\true code} {\false code}

```

Compares the definition of two *<control sequences>* and is logically `true` if they are the same, *i.e.* if they have exactly the same definition when examined with `\cs_show:N`.

```

\cs_if_exist_p:N *
\cs_if_exist_p:c *
\cs_if_exist:NTF *
\cs_if_exist:cTF *

```

```

\cs_if_exist_p:N <control sequence>
\cs_if_exist:NNTF <control sequence> {\true code} {\false code}

```

Tests whether the *<control sequence>* is currently defined (whether as a function or another control sequence type). Any definition of *<control sequence>* other than `\relax` evaluates as `true`.

```

\cs_if_free_p:N *
\cs_if_free_p:c *
\cs_if_free:NTF *
\cs_if_free:cTF *

```

```

\cs_if_free_p:N <control sequence>
\cs_if_free:NNTF <control sequence> {\true code} {\false code}

```

Tests whether the *<control sequence>* is currently free to be defined. This test is `false` if the *<control sequence>* currently exists (as defined by `\cs_if_exist:N`).

6.2 Primitive conditionals

The ε -T_EX engine itself provides many different conditionals. Some expand whatever comes after them and others don’t. Hence the names for these underlying functions often contains a `:w` part but higher level functions are often available. See for instance `\int_compare_p:nNn` which is a wrapper for `\if_int_compare:w`.

Certain conditionals deal with specific data types like boxes and fonts and are described there. The ones described below are either the universal conditionals or deal with control sequences. We prefix primitive conditionals with `\if_`.

<code>\if_true:</code>	★	<code>\if_true: <true code> \else: <false code> \fi:</code>
<code>\if_false:</code>	★	<code>\if_false: <true code> \else: <false code> \fi:</code>
<code>\else:</code>	★	<code>\reverse_if:N <primitive conditional></code>
<code>\fi:</code>	★	<code>\if_true:</code> always executes <i><true code></i> , while <code>\if_false:</code> always executes <i><false code></i> .
<code>\reverse_if:N</code>	★	<code>\reverse_if:N</code> reverses any two-way primitive conditional. <code>\else:</code> and <code>\fi:</code> delimit the branches of the conditional. The function <code>\or:</code> is documented in <code>l3int</code> and used in case switches.

TeXhackers note: These are equivalent to their corresponding TeX primitive conditionals; `\reverse_if:N` is ε -TeX's `\unless`.

<code>\if_meaning:w</code>	★	<code>\if_meaning:w <arg₁₂</code>
----------------------------	---	---

`\if_meaning:w` executes *<true code>* when *<arg_{1 and *<arg_{2 are the same, otherwise it executes *<false code>*. *<arg_{1 and *<arg_{2 could be functions, variables, tokens; in all cases the *unexpanded* definitions are compared.}*}*}*}*

TeXhackers note: This is TeX's `\ifx`.

<code>\if:w</code>	★	<code>\if:w <token₁₂</code>
<code>\if_charcode:w</code>	★	<code>\if_catcode:w <token₁₂</code>
<code>\if_catcode:w</code>	★	These conditionals expand any following tokens until two unexpandable tokens are left. If you wish to prevent this expansion, prefix the token in question with <code>\exp_not:N</code> . <code>\if_catcode:w</code> tests if the category codes of the two tokens are the same whereas <code>\if:w</code> tests if the character codes are identical. <code>\if_charcode:w</code> is an alternative name for <code>\if:w</code> .

<code>\if_cs_exist:N</code>	★	<code>\if_cs_exist:N <cs> <true code> \else: <false code> \fi:</code>
<code>\if_cs_exist:w</code>	★	<code>\if_cs_exist:w <tokens> \cs_end: <true code> \else: <false code> \fi:</code>

Check if *<cs>* appears in the hash table or if the control sequence that can be formed from *<tokens>* appears in the hash table. The latter function does not turn the control sequence in question into `\scan_stop:!` This can be useful when dealing with control sequences which cannot be entered as a single token.

<code>\if_mode_horizontal:</code>	★	<code>\if_mode_horizontal: <true code> \else: <false code> \fi:</code>
<code>\if_mode_vertical:</code>	★	Execute <i><true code></i> if currently in horizontal mode, otherwise execute <i><false code></i> . Similar for the other functions.
<code>\if_mode_math:</code>	★	
<code>\if_mode_inner:</code>	★	

7 Starting a paragraph

`\mode_leave_vertical:`

New: 2017-07-04

`\mode_leave_vertical:`

Ensures that `TEX` is not in vertical (inter-paragraph) mode. In horizontal or math mode this command has no effect, in vertical mode it switches to horizontal mode, and inserts a box of width `\parindent`, followed by the `\everypar` token list.

T_EXhackers note: This results in the contents of the `\everypar` token register being inserted, after `\mode_leave_vertical:` is complete. Notice that in contrast to the L^AT_EX 2_ε `\leavevmode` approach, no box is used by the method implemented here.

8 Debugging support

`\debug_on:n`
`\debug_off:n`

New: 2017-07-16
Updated: 2017-08-02

`\debug_on:n { <comma-separated list> }`
`\debug_off:n { <comma-separated list> }`

Turn on and off within a group various debugging code, some of which is also available as `expl3` load-time options. The items that can be used in the *<list>* are

- **check-declarations** that checks all `expl3` variables used were previously declared and that local/global variables (based on their name or on their first assignment) are only locally/globally assigned;
- **check-expressions** that checks integer, dimension, skip, and muskip expressions are not terminated prematurely;
- **deprecation** that makes soon-to-be-deprecated commands produce errors;
- **log-functions** that logs function definitions;
- **all** that does all of the above.

Providing these as switches rather than options allows testing code even if it relies on other packages: load all other packages, call `\debug_on:n`, and load the code that one is interested in testing. These functions can only be used in L^AT_EX 2_ε package mode loaded with `enable-debug` or another option implying it.

`\debug_suspend:`
`\debug_resume:`

New: 2017-11-28

`\debug_suspend: ... \debug_resume:`

Suppress (locally) errors and logging from `debug` commands, except for the **deprecation** errors or warnings. These pairs of commands can be nested. This can be used around pieces of code that are known to fail checks, if such failures should be ignored. See for instance `l3coffins`.

Part V

The l3expan package

Argument expansion

This module provides generic methods for expanding T_EX arguments in a systematic manner. The functions in this module all have prefix `exp`.

Not all possible variations are implemented for every base function. Instead only those that are used within the L^AT_EX3 kernel or otherwise seem to be of general interest are implemented. Consult the module description to find out which functions are actually defined. The next section explains how to define missing variants.

1 Defining new variants

The definition of variant forms for base functions may be necessary when writing new functions or when applying a kernel function in a situation that we haven't thought of before.

Internally preprocessing of arguments is done with functions of the form `\exp_....`. They all look alike, an example would be `\exp_args:NNo`. This function has three arguments, the first and the second are a single tokens, while the third argument should be given in braces. Applying `\exp_args:NNo` expands the content of third argument once before any expansion of the first and second arguments. If `\seq_gpush:No` was not defined it could be coded in the following way:

```
\exp_args:NNo \seq_gpush:Nn
  \g_file_name_stack
  { \l_tmpa_tl }
```

In other words, the first argument to `\exp_args:NNo` is the base function and the other arguments are preprocessed and then passed to this base function. In the example the first argument to the base function should be a single token which is left unchanged while the second argument is expanded once. From this example we can also see how the variants are defined. They just expand into the appropriate `\exp_` function followed by the desired base function, *e.g.*

```
\cs_generate_variant:Nn \seq_gpush:Nn { No }
```

results in the definition of `\seq_gpush:No`

```
\cs_new:Npn \seq_gpush:No { \exp_args:NNo \seq_gpush:Nn }
```

Providing variants in this way in style files is safe as the `\cs_generate_variant:Nn` function will only create new definitions if there is not already one available. Therefore adding such definition to later releases of the kernel will not make such style files obsolete.

The steps above may be automated by using the function `\cs_generate_variant:Nn`, described next.

2 Methods for defining variants

We recall the set of available argument specifiers.

- `N` is used for single-token arguments while `c` constructs a control sequence from its name and passes it to a parent function as an `N`-type argument.
- Many argument types extract or expand some tokens and provide it as an `n`-type argument, namely a braced multiple-token argument: `V` extracts the value of a variable, `v` extracts the value from the name of a variable, `n` uses the argument as it is, `o` expands once, `f` expands fully the front of the token list, `e` and `x` expand fully all tokens (differences are explained later).
- A few odd argument types remain: `T` and `F` for conditional processing, otherwise identical to `n`-type arguments, `p` for the parameter text in definitions, `w` for arguments with a specific syntax, and `D` to denote primitives that should not be used directly.

`\cs_generate_variant:Nn`
`\cs_generate_variant:cn`

Updated: 2017-11-28

`\cs_generate_variant:Nn` \langle parent control sequence \rangle $\{$ \langle variant argument specifiers \rangle $\}$

This function is used to define argument-specifier variants of the \langle parent control sequence \rangle for L^AT_EX3 code-level macros. The \langle parent control sequence \rangle is first separated into the \langle base name \rangle and \langle original argument specifier \rangle . The comma-separated list of \langle variant argument specifiers \rangle is then used to define variants of the \langle original argument specifier \rangle if these are not already defined. For each \langle variant \rangle given, a function is created that expands its arguments as detailed and passes them to the \langle parent control sequence \rangle . So for example

```
\cs_set:Npn \foo:Nn #1#2 { code here }
\cs_generate_variant:Nn \foo:Nn { c }
```

creates a new function `\foo:cn` which expands its first argument into a control sequence name and passes the result to `\foo:Nn`. Similarly

```
\cs_generate_variant:Nn \foo:Nn { NV , cV }
```

generates the functions `\foo:NV` and `\foo:cV` in the same way. The `\cs_generate_variant:Nn` function can only be applied if the \langle parent control sequence \rangle is already defined. If the \langle parent control sequence \rangle is protected or if the \langle variant \rangle involves any **x** argument, then the \langle variant control sequence \rangle is also protected. The \langle variant \rangle is created globally, as is any `\exp_args:N` \langle variant \rangle function needed to carry out the expansion.

Only **n** and **N** arguments can be changed to other types. The only allowed changes are

- **c** variant of an **N** parent;
- **o**, **V**, **v**, **f**, **e**, or **x** variant of an **n** parent;
- **N**, **n**, **T**, **F**, or **p** argument unchanged.

This means the \langle parent \rangle of a \langle variant \rangle form is always unambiguous, even in cases where both an **n**-type parent and an **N**-type parent exist, such as for `\tl_count:n` and `\tl_count:N`.

For backward compatibility it is currently possible to make **n**, **o**, **V**, **v**, **f**, **e**, or **x**-type variants of an **N**-type argument or **N** or **c**-type variants of an **n**-type argument. Both are deprecated. The first because passing more than one token to an **N**-type argument will typically break the parent function's code. The second because programmers who use that most often want to access the value of a variable given its name, hence should use a **V**-type or **v**-type variant instead of **c**-type. In those cases, using the lower-level `\exp_args:No` or `\exp_args:Nc` functions explicitly is preferred to defining confusing variants.

3 Introducing the variants

The **V** type returns the value of a register, which can be one of **tl**, **clist**, **int**, **skip**, **dim**, **muskip**, or built-in T_EX registers. The **v** type is the same except it first creates a control sequence out of its argument before returning the value.

In general, the programmer should not need to be concerned with expansion control. When simply using the content of a variable, functions with a **V** specifier should be used. For those referred to by (cs)name, the **v** specifier is available for the same purpose. Only

when specific expansion steps are needed, such as when using delimited arguments, should the lower-level functions with `o` specifiers be employed.

The `e` type expands all tokens fully, starting from the first. More precisely the expansion is identical to that of T_EX’s `\message` (in particular `#` needs not be doubled). It was added in May 2018. In recent enough engines (starting around 2019) it relies on the primitive `\expanded` hence is fast. In older engines it is very much slower. As a result it should only be used in performance critical code if typical users will have a recent installation of the T_EX ecosystem.

The `x` type expands all tokens fully, starting from the first. In contrast to `e`, all macro parameter characters `#` must be doubled, and omitting this leads to low-level errors. In addition this type of expansion is not expandable, namely functions that have `x` in their signature do not themselves expand when appearing inside `x` or `e` expansion.

The `f` type is so special that it deserves an example. It is typically used in contexts where only expandable commands are allowed. Then `x`-expansion cannot be used, and `f`-expansion provides an alternative that expands the front of the token list as much as can be done in such contexts. For instance, say that we want to evaluate the integer expression `3 + 4` and pass the result 7 as an argument to an expandable function `\example:n`. For this, one should define a variant using `\cs_generate_variant:Nn \example:n { f }`, then do

```
\example:f { \int_eval:n { 3 + 4 } }
```

Note that `x`-expansion would also expand `\int_eval:n` fully to its result 7, but the variant `\example:x` cannot be expandable. Note also that `o`-expansion would not expand `\int_eval:n` fully to its result since that function requires several expansions. Besides the fact that `x`-expansion is protected rather than expandable, another difference between `f`-expansion and `x`-expansion is that `f`-expansion expands tokens from the beginning and stops as soon as a non-expandable token is encountered, while `x`-expansion continues expanding further tokens. Thus, for instance

```
\example:f { \int_eval:n { 1 + 2 } , \int_eval:n { 3 + 4 } }
```

results in the call

```
\example:n { 3 , \int_eval:n { 3 + 4 } }
```

while using `\example:x` or `\example:e` instead results in

```
\example:n { 3 , 7 }
```

at the cost of being protected (for `x` type) or very much slower in old engines (for `e` type). If you use `f` type expansion in conditional processing then you should stick to using TF type functions only as the expansion does not finish any `\if... \fi`: itself!

It is important to note that both `f`- and `o`-type expansion are concerned with the expansion of tokens from left to right in their arguments. In particular, `o`-type expansion applies to the first *token* in the argument it receives: it is conceptually similar to

```
\exp_after:wN <base function> \exp_after:wN { <argument> }
```

At the same time, `f`-type expansion stops at the *first* non-expandable token. This means for example that both

```
\tl_set:No \l_tmpa_tl { { \g_tmpb_tl } }
```

and

```
\tl_set:Nf \l_tmpa_tl { { \g_tmpb_tl } }
```

leave `\g_tmpb_tl` unchanged: `{` is the first token in the argument and is non-expandable. It is usually best to keep the following in mind when using variant forms.

- Variants with `x`-type arguments (that are fully expanded before being passed to the `n`-type base function) are never expandable even when the base function is. Such variants cannot work correctly in arguments that are themselves subject to expansion. Consider using `f` or `e` expansion.
- In contrast, `e` expansion (full expansion, almost like `x` except for the treatment of `#`) does not prevent variants from being expandable (if the base function is). The drawback is that `e` expansion is very much slower in old engines (before 2019). Consider using `f` expansion if that type of expansion is sufficient to perform the required expansion, or `x` expansion if the variant will not itself need to be expandable.
- Finally `f` expansion only expands the front of the token list, stopping at the first non-expandable token. This may fail to fully expand the argument.

When speed is essential (for functions that do very little work and whose variants are used numerous times in a document) the following considerations apply because internal functions for argument expansion come in two flavours, some faster than others.

- Arguments that might need expansion should come first in the list of arguments.
- Arguments that should consist of single tokens `N`, `c`, `V`, or `v` should come first among these.
- Arguments that appear after the first multi-token argument `n`, `f`, `e`, or `o` require slightly slower special processing to be expanded. Therefore it is best to use the optimized functions, namely those that contain only `N`, `c`, `V`, and `v`, and, in the last position, `o`, `f`, `e`, with possible trailing `N` or `n` or `T` or `F`, which are not expanded. Any `x`-type argument causes slightly slower processing.

4 Manipulating the first argument

These functions are described in detail: expansion of multiple tokens follows the same rules but is described in a shorter fashion.

```
\exp_args:Nc ★ \exp_args:Nc <function> {<tokens>}  
\exp_args:cc ★
```

This function absorbs two arguments (the `<function>` name and the `<tokens>`). The `<tokens>` are expanded until only characters remain, and are then turned into a control sequence. The result is inserted into the input stream *after* reinsertion of the `<function>`. Thus the `<function>` may take more than one argument: all others are left unchanged.

The `:cc` variant constructs the `<function>` name in the same manner as described for the `<tokens>`.

T_EXhackers note: Protected macros that appear in a `c`-type argument are expanded despite being protected; `\exp_not:n` also has no effect. An internal error occurs if non-characters or active characters remain after full expansion, as the conversion to a control sequence is not possible.

`\exp_args:No` ★ `\exp_args:No` $\langle function \rangle$ $\{\langle tokens \rangle\}$...

This function absorbs two arguments (the $\langle function \rangle$ name and the $\langle tokens \rangle$). The $\langle tokens \rangle$ are expanded once, and the result is inserted in braces into the input stream *after* reinsertion of the $\langle function \rangle$. Thus the $\langle function \rangle$ may take more than one argument: all others are left unchanged.

`\exp_args:NV` ★ `\exp_args:NV` $\langle function \rangle$ $\langle variable \rangle$

This function absorbs two arguments (the names of the $\langle function \rangle$ and the $\langle variable \rangle$). The content of the $\langle variable \rangle$ are recovered and placed inside braces into the input stream *after* reinsertion of the $\langle function \rangle$. Thus the $\langle function \rangle$ may take more than one argument: all others are left unchanged.

`\exp_args:Nv` ★ `\exp_args:Nv` $\langle function \rangle$ $\{\langle tokens \rangle\}$

This function absorbs two arguments (the $\langle function \rangle$ name and the $\langle tokens \rangle$). The $\langle tokens \rangle$ are expanded until only characters remain, and are then turned into a control sequence. This control sequence should be the name of a $\langle variable \rangle$. The content of the $\langle variable \rangle$ are recovered and placed inside braces into the input stream *after* reinsertion of the $\langle function \rangle$. Thus the $\langle function \rangle$ may take more than one argument: all others are left unchanged.

TeXhackers note: Protected macros that appear in a v-type argument are expanded despite being protected; `\exp_not:n` also has no effect. An internal error occurs if non-characters or active characters remain after full expansion, as the conversion to a control sequence is not possible.

`\exp_args:Ne` ★ `\exp_args:Ne` $\langle function \rangle$ $\{\langle tokens \rangle\}$

New: 2018-05-15

This function absorbs two arguments (the $\langle function \rangle$ name and the $\langle tokens \rangle$) and exhaustively expands the $\langle tokens \rangle$. The result is inserted in braces into the input stream *after* reinsertion of the $\langle function \rangle$. Thus the $\langle function \rangle$ may take more than one argument: all others are left unchanged.

TeXhackers note: This relies on the `\expanded` primitive when available (in LuaTeX and starting around 2019 in other engines). Otherwise it uses some fall-back code that is very much slower. As a result it should only be used in performance-critical code if typical users have a recent installation of the TeX ecosystem.

`\exp_args:Nf` ★ `\exp_args:Nf` $\langle function \rangle$ $\{\langle tokens \rangle\}$

This function absorbs two arguments (the $\langle function \rangle$ name and the $\langle tokens \rangle$). The $\langle tokens \rangle$ are fully expanded until the first non-expandable token is found (if that is a space it is removed), and the result is inserted in braces into the input stream *after* reinsertion of the $\langle function \rangle$. Thus the $\langle function \rangle$ may take more than one argument: all others are left unchanged.

<code>\exp_args:Nx</code>	<code>\exp_args:Nx</code>	$\langle function \rangle$	$\{\langle tokens \rangle\}$
---------------------------	---------------------------	----------------------------	------------------------------

This function absorbs two arguments (the $\langle function \rangle$ name and the $\langle tokens \rangle$) and exhaustively expands the $\langle tokens \rangle$. The result is inserted in braces into the input stream *after* reinsertion of the $\langle function \rangle$. Thus the $\langle function \rangle$ may take more than one argument: all others are left unchanged.

5 Manipulating two arguments

<code>\exp_args:NNc</code>	<code>\exp_args:NNc</code>	$\langle token_1 \rangle$	$\langle token_2 \rangle$	$\{\langle tokens \rangle\}$
----------------------------	----------------------------	---------------------------	---------------------------	------------------------------

These optimized functions absorb three arguments and expand the second and third as detailed by their argument specifier. The first argument of the function is then the next item on the input stream, followed by the expansion of the second and third arguments.

`\exp_args:NNv` *
`\exp_args:NNe` *
`\exp_args:NNf` *
`\exp_args:Ncc` *
`\exp_args:Nco` *
`\exp_args:NcV` *
`\exp_args:Ncv` *
`\exp_args:Ncf` *
`\exp_args:NVV` *

Updated: 2018-05-15

<code>\exp_args:Nnc</code>	<code>\exp_args:Noo</code>	$\langle token \rangle$	$\{\langle tokens_1 \rangle\}$	$\{\langle tokens_2 \rangle\}$
----------------------------	----------------------------	-------------------------	--------------------------------	--------------------------------

These functions absorb three arguments and expand the second and third as detailed by their argument specifier. The first argument of the function is then the next item on the input stream, followed by the expansion of the second and third arguments. These functions need slower processing.

`\exp_args:Noc` *
`\exp_args:Noo` *
`\exp_args:Nof` *
`\exp_args:NVo` *
`\exp_args:Nfo` *
`\exp_args:Nff` *
`\exp_args:Nee` *

Updated: 2018-05-15

<code>\exp_args:NNx</code>	<code>\exp_args:NNx</code>	$\langle token_1 \rangle$	$\langle token_2 \rangle$	$\{\langle tokens \rangle\}$
----------------------------	----------------------------	---------------------------	---------------------------	------------------------------

These functions absorb three arguments and expand the second and third as detailed by their argument specifier. The first argument of the function is then the next item on the input stream, followed by the expansion of the second and third arguments. These functions are not expandable due to their x-type argument.

`\exp_args:Ncx` *
`\exp_args:Nnx` *
`\exp_args:Nox` *
`\exp_args:Nxo` *
`\exp_args:Nxx` *

6 Manipulating three arguments

<code>\exp_args:NNNo</code>	*	<code>\exp_args:NNNo</code>	$\langle token_1 \rangle$	$\langle token_2 \rangle$	$\langle token_3 \rangle$	$\{\langle tokens \rangle\}$
<code>\exp_args:NNNV</code>	*	These optimized functions absorb four arguments and expand the second, third and fourth as detailed by their argument specifier. The first argument of the function is then the next item on the input stream, followed by the expansion of the second argument, <i>etc.</i>				
<code>\exp_args:NNNv</code>	*					
<code>\exp_args:Nccc</code>	*					
<code>\exp_args:NcNc</code>	*					
<code>\exp_args:NcNo</code>	*					
<code>\exp_args:Ncco</code>	*					

<code>\exp_args:NNcf</code>	*	<code>\exp_args:NNoo</code>	$\langle token_1 \rangle$	$\langle token_2 \rangle$	$\{\langle token_3 \rangle\}$	$\{\langle tokens \rangle\}$
<code>\exp_args:NNno</code>	*	These functions absorb four arguments and expand the second, third and fourth as detailed by their argument specifier. The first argument of the function is then the next item on the input stream, followed by the expansion of the second argument, <i>etc.</i> These functions need slower processing.				
<code>\exp_args:NNnV</code>	*					
<code>\exp_args:NNoo</code>	*					
<code>\exp_args:NNVV</code>	*					
<code>\exp_args:Ncno</code>	*					
<code>\exp_args:NcnV</code>	*					
<code>\exp_args:Ncoo</code>	*					
<code>\exp_args:NcVV</code>	*					
<code>\exp_args:Nnnc</code>	*					
<code>\exp_args:Nnno</code>	*					
<code>\exp_args:Nnnf</code>	*					
<code>\exp_args:Nnff</code>	*					
<code>\exp_args:Nooo</code>	*					
<code>\exp_args:Noof</code>	*					
<code>\exp_args:Nffo</code>	*					
<code>\exp_args:Neee</code>	*					

<code>\exp_args:NNNx</code>	<code>\exp_args:NNnx</code>	$\langle token_1 \rangle$	$\langle token_2 \rangle$	$\{\langle tokens_1 \rangle\}$	$\{\langle tokens_2 \rangle\}$
<code>\exp_args:NNox</code>	These functions absorb four arguments and expand the second, third and fourth as detailed by their argument specifier. The first argument of the function is then the next item on the input stream, followed by the expansion of the second argument, <i>etc.</i>				
<code>\exp_args:Nccx</code>					
<code>\exp_args:Ncnx</code>					
<code>\exp_args:Nnnx</code>					
<code>\exp_args:Nnox</code>					
<code>\exp_args:Noox</code>					

New: 2015-08-12

7 Unbraced expansion

```

\exp_last_unbraced:No  *
\exp_last_unbraced:Nv  *
\exp_last_unbraced:Nf  *
\exp_last_unbraced:NNo *
\exp_last_unbraced:NNV *
\exp_last_unbraced:NNf *
\exp_last_unbraced:Nco *
\exp_last_unbraced:NcV *
\exp_last_unbraced:Nno *
\exp_last_unbraced:Noo *
\exp_last_unbraced:Nfo *
\exp_last_unbraced:NNNo *
\exp_last_unbraced:NNNV *
\exp_last_unbraced:NNNf *
\exp_last_unbraced:NnNo *
\exp_last_unbraced:NNNNo *
\exp_last_unbraced:NNNNf *

```

Updated: 2018-05-15

```
\exp_last_unbraced:Nno <token> {\tokens_1} {\tokens_2}
```

These functions absorb the number of arguments given by their specification, carry out the expansion indicated and leave the results in the input stream, with the last argument not surrounded by the usual braces. Of these, the :Nno, :Noo, :Nfo and :NnNo variants need slower processing.

T_EXhackers note: As an optimization, the last argument is unbraced by some of those functions before expansion. This can cause problems if the argument is empty: for instance, `\exp_last_unbraced:Nf \foo_bar:w { } \q_stop` leads to an infinite loop, as the quark is f-expanded.

```
\exp_last_unbraced:Nx \exp_last_unbraced:Nx <function> {\tokens}
```

This function fully expands the `<tokens>` and leaves the result in the input stream after reinsertion of the `<function>`. This function is not expandable.

```
\exp_last_two_unbraced:Noo * \exp_last_two_unbraced:Noo <token> {\tokens_1} {\tokens_2}
```

This function absorbs three arguments and expands the second and third once. The first argument of the function is then the next item on the input stream, followed by the expansion of the second and third arguments, which are not wrapped in braces. This function needs special (slower) processing.

```
\exp_after:wN * \exp_after:wN <token_1> <token_2>
```

Carries out a single expansion of `<token_2>` (which may consume arguments) prior to the expansion of `<token_1>`. If `<token_2>` has no expansion (for example, if it is a character) then it is left unchanged. It is important to notice that `<token_1>` may be *any* single token, including group-opening and -closing tokens (`{` or `}` assuming normal T_EX category codes). Unless specifically required this should be avoided: expansion should be carried out using an appropriate argument specifier variant or the appropriate `\exp_arg:N` function.

T_EXhackers note: This is the T_EX primitive `\expandafter` renamed.

8 Preventing expansion

Despite the fact that the following functions are all about preventing expansion, they're designed to be used in an expandable context and hence are all marked as being 'expand-

able' since they themselves disappear after the expansion has completed.

`\exp_not:N` ★ `\exp_not:N` $\langle token \rangle$

Prevents expansion of the $\langle token \rangle$ in a context where it would otherwise be expanded, for example an **x**-type argument or the first token in an **o** or **e** or **f** argument.

T_EXhackers note: This is the T_EX `\noexpand` primitive. It only prevents expansion. At the beginning of an **f**-type argument, a space $\langle token \rangle$ is removed even if it appears as `\exp_not:N \c_space_token`. In an **x**-expanding definition (`\cs_new:Npx`), a macro parameter introduces an argument even if it appears as `\exp_not:N # 1`. This differs from `\exp_not:n`.

`\exp_not:c` ★ `\exp_not:c` $\{\langle tokens \rangle\}$

Expands the $\langle tokens \rangle$ until only characters remain, and then converts this into a control sequence. Further expansion of this control sequence is then inhibited using `\exp_not:N`.

T_EXhackers note: Protected macros that appear in a **c**-type argument are expanded despite being protected; `\exp_not:n` also has no effect. An internal error occurs if non-characters or active characters remain after full expansion, as the conversion to a control sequence is not possible.

`\exp_not:n` ★ `\exp_not:n` $\{\langle tokens \rangle\}$

Prevents expansion of the $\langle tokens \rangle$ in an **e** or **x**-type argument. In all other cases the $\langle tokens \rangle$ continue to be expanded, for example in the input stream or in other types of arguments such as **c**, **f**, **v**. The argument of `\exp_not:n` *must* be surrounded by braces.

T_EXhackers note: This is the ε -T_EX `\unexpanded` primitive. In an **x**-expanding definition (`\cs_new:Npx`), `\exp_not:n {#1}` is equivalent to `##1` rather than to `#1`, namely it inserts the two characters `#` and `1`. In an **e**-type argument `\exp_not:n {#}` is equivalent to `#`, namely it inserts the character `#`.

`\exp_not:o` ★ `\exp_not:o` $\{\langle tokens \rangle\}$

Expands the $\langle tokens \rangle$ once, then prevents any further expansion in **x**-type or **e**-type arguments using `\exp_not:n`.

`\exp_not:V` ★ `\exp_not:V` $\langle variable \rangle$

Recovers the content of the $\langle variable \rangle$, then prevents expansion of this material in **x**-type or **e**-type arguments using `\exp_not:n`.

<code>\exp_not:v</code>	<code>\exp_not:v {<tokens>}</code>
-------------------------	--

Expands the *<tokens>* until only characters remains, and then converts this into a control sequence which should be a *<variable>* name. The content of the *<variable>* is recovered, and further expansion in *x*-type or *e*-type arguments is prevented using `\exp_not:n`.

TeXhackers note: Protected macros that appear in a *v*-type argument are expanded despite being protected; `\exp_not:n` also has no effect. An internal error occurs if non-characters or active characters remain after full expansion, as the conversion to a control sequence is not possible.

<code>\exp_not:e</code>	<code>\exp_not:e {<tokens>}</code>
-------------------------	--

Expands *<tokens>* exhaustively, then protects the result of the expansion (including any tokens which were not expanded) from further expansion in *e* or *x*-type arguments using `\exp_not:n`. This is very rarely useful but is provided for consistency.

<code>\exp_not:f</code>	<code>\exp_not:f {<tokens>}</code>
-------------------------	--

Expands *<tokens>* fully until the first unexpandable token is found (if it is a space it is removed). Expansion then stops, and the result of the expansion (including any tokens which were not expanded) is protected from further expansion in *x*-type or *e*-type arguments using `\exp_not:n`.

<code>\exp_stop_f:</code>	<code>\foo_bar:f { <tokens> \exp_stop_f: <more tokens> }</code>
---------------------------	---

Updated: 2011-06-03

This function terminates an *f*-type expansion. Thus if a function `\foo_bar:f` starts an *f*-type expansion and all of *<tokens>* are expandable `\exp_stop_f:` terminates the expansion of tokens even if *<more tokens>* are also expandable. The function itself is an implicit space token. Inside an *x*-type expansion, it retains its form, but when typeset it produces the underlying space (␣).

9 Controlled expansion

The `expl3` language makes all efforts to hide the complexity of TeX expansion from the programmer by providing concepts that evaluate/expand arguments of functions prior to calling the “base” functions. Thus, instead of using many `\expandafter` calls and other trickery it is usually a matter of choosing the right variant of a function to achieve a desired result.

Of course, deep down TeX is using expansion as always and there are cases where a programmer needs to control that expansion directly; typical situations are basic data manipulation tools. This section documents the functions for that level. These commands are used throughout the kernel code, but we hope that outside the kernel there will be little need to resort to them. Instead the argument manipulation methods document above should usually be sufficient.

While `\exp_after:wN` expands one token (out of order) it is sometimes necessary to expand several tokens in one go. The next set of commands provide this functionality. Be aware that it is absolutely required that the programmer has full control over the tokens to be expanded, i.e., it is not possible to use these functions to expand unknown input as part of *<expandable-tokens>* as that will break badly if unexpandable tokens are encountered in that place!

<code>\exp:w</code>	★	<code>\exp:w <expandable tokens> \exp_end:</code>
<code>\exp_end:</code>	★	Expands <code><expandable-tokens></code> until reaching <code>\exp_end:</code> at which point expansion stops. The full expansion of <code><expandable tokens></code> has to be empty. If any token in <code><expandable tokens></code> or any token generated by expanding the tokens therein is not expandable the expansion will end prematurely and as a result <code>\exp_end:</code> will be misinterpreted later on. ⁴
New: 2015-08-23		

In typical use cases the `\exp_end:` is hidden somewhere in the replacement text of `<expandable-tokens>` rather than being on the same expansion level than `\exp:w`, e.g., you may see code such as

```
\exp:w \@@_case:NnTF #1 {#2} { } { }
```

where somewhere during the expansion of `\@@_case:NnTF` the `\exp_end:` gets generated.

T_EXhackers note: The current implementation uses `\romannumeral` hence ignores space tokens and explicit signs + and - in the expansion of the `<expandable tokens>`, but this should not be relied upon.

<code>\exp:w</code>	★	<code>\exp:w <expandable-tokens> \exp_end_continue_f:w <further-tokens></code>
<code>\exp_end_continue_f:w</code>	★	Expands <code><expandable-tokens></code> until reaching <code>\exp_end_continue_f:w</code> at which point expansion continues as an f-type expansion expanding <code><further-tokens></code> until an unexpandable token is encountered (or the f-type expansion is explicitly terminated by <code>\exp_stop_f:</code>). As with all f-type expansions a space ending the expansion gets removed. The full expansion of <code><expandable-tokens></code> has to be empty. If any token in <code><expandable-tokens></code> or any token generated by expanding the tokens therein is not expandable the expansion will end prematurely and as a result <code>\exp_end_continue_f:w</code> will be misinterpreted later on. ⁵
New: 2015-08-23		

The full expansion of `<expandable-tokens>` has to be empty. If any token in `<expandable-tokens>` or any token generated by expanding the tokens therein is not expandable the expansion will end prematurely and as a result `\exp_end_continue_f:w` will be misinterpreted later on.⁵

In typical use cases `<expandable-tokens>` contains no tokens at all, e.g., you will see code such as

```
\exp_after:wN { \exp:w \exp_end_continue_f:w #2 }
```

where the `\exp_after:wN` triggers an f-expansion of the tokens in #2. For technical reasons this has to happen using two tokens (if they would be hidden inside another command `\exp_after:wN` would only expand the command but not trigger any additional f-expansion).

You might wonder why there are two different approaches available, after all the effect of

```
\exp:w <expandable-tokens> \exp_end:
```

can be alternatively achieved through an f-type expansion by using `\exp_stop_f:`, i.e.

```
\exp:w \exp_end_continue_f:w <expandable-tokens> \exp_stop_f:
```

The reason is simply that the first approach is slightly faster (one less token to parse and less expansion internally) so in places where such performance really matters and where we want to explicitly stop the expansion at a defined point the first form is preferable.

⁴Due to the implementation you might get the character in position 0 in the current font (typically “”) in the output without any error message!

⁵In this particular case you may get a character into the output as well as an error message.

<code>\exp:w</code>	★
<code>\exp_end_continue_f:nw</code>	★

New: 2015-08-23

`\exp:w` *<expandable-tokens>* `\exp_end_continue_f:nw` *<further-tokens>*

The difference to `\exp_end_continue_f:w` is that we first we pick up an argument which is then returned to the input stream. If *<further-tokens>* starts with space tokens then these space tokens are removed while searching for the argument. If it starts with a brace group then the braces are removed. Thus such spaces or braces will not terminate the f-type expansion.

10 Internal functions

`\::n` `\cs_new:Npn \exp_args:Ncof { \::c \::o \::f \::: }`

`\::N` Internal forms for the base expansion types. These names do *not* conform to the general \LaTeX 3 approach as this makes them more readily visible in the log and so forth. They should not be used outside this module.

`\::c`
`\::o`
`\::e`
`\::f`
`\::x`
`\::v`
`\::V`
`\:::`

`\::o_unbraced` `\cs_new:Npn \exp_last_unbraced:Nno { \::n \::o_unbraced \::: }`

`\::e_unbraced` Internal forms for the expansion types which leave the terminal argument unbraced. These names do *not* conform to the general \LaTeX 3 approach as this makes them more readily visible in the log and so forth. They should not be used outside this module.

`\::f_unbraced`
`\::x_unbraced`
`\::v_unbraced`
`\::V_unbraced`

Part VI

The l3quark package

Quarks

Two special types of constants in L^AT_EX3 are “quarks” and “scan marks”. By convention all constants of type quark start out with `\q_`, and scan marks start with `\s_`.

1 Quarks

Quarks are control sequences that expand to themselves and should therefore *never* be executed directly in the code. This would result in an endless loop!

They are meant to be used as delimiter in weird functions, the most common use case being the ‘stop token’ (*i.e.* `\q_stop`). For example, when writing a macro to parse a user-defined date

```
\date_parse:n {19/June/1981}
```

one might write a command such as

```
\cs_new:Npn \date_parse:n #1 { \date_parse_aux:w #1 \q_stop }
\cs_new:Npn \date_parse_aux:w #1 / #2 / #3 \q_stop
{ <do something with the date> }
```

Quarks are sometimes also used as error return values for functions that receive erroneous input. For example, in the function `\prop_get:NnN` to retrieve a value stored in some key of a property list, if the key does not exist then the return value is the quark `\q_no_value`. As mentioned above, such quarks are extremely fragile and it is imperative when using such functions that code is carefully written to check for pathological cases to avoid leakage of a quark into an uncontrolled environment.

Quarks also permit the following ingenious trick when parsing tokens: when you pick up a token in a temporary variable and you want to know whether you have picked up a particular quark, all you have to do is compare the temporary variable to the quark using `\tl_if_eq:NNTF`. A set of special quark testing functions is set up below. All the quark testing functions are expandable although the ones testing only single tokens are much faster. An example of the quark testing functions and their use in recursion can be seen in the implementation of `\clist_map_function:NN`.

2 Defining quarks

`\quark_new:N`

`\quark_new:N <quark>`

Creates a new `<quark>` which expands only to `<quark>`. The `<quark>` is defined globally, and an error message is raised if the name was already taken.

`\q_stop`

Used as a marker for delimited arguments, such as

```
\cs_set:Npn \tmp:w #1#2 \q_stop {#1}
```

<u><u>\q_mark</u></u>	Used as a marker for delimited arguments when <code>\q_stop</code> is already in use.
<u><u>\q_nil</u></u>	Quark to mark a null value in structured variables or functions. Used as an end delimiter when this may itself need to be tested (in contrast to <code>\q_stop</code> , which is only ever used as a delimiter).
<u><u>\q_no_value</u></u>	A canonical value for a missing value, when one is requested from a data structure. This is therefore used as a “return” value by functions such as <code>\prop_get:NnN</code> if there is no data to return.

3 Quark tests

The method used to define quarks means that the single token (N) tests are faster than the multi-token (n) tests. The latter should therefore only be used when the argument can definitely take more than a single token.

<u><u>\quark_if_nil_p:N</u></u> *	<u><u>\quark_if_nil:N</u></u> <i><token></i>
<u><u>\quark_if_nil:N</u></u> <i>TF</i> *	<u><u>\quark_if_nil:N</u></u> <i>TF</i> <i><token></i> <i>{<true code>}</i> <i>{<false code>}</i>
	Tests if the <i><token></i> is equal to <code>\q_nil</code> .
<u><u>\quark_if_nil_p:n</u></u> *	<u><u>\quark_if_nil_p:n</u></u> <i>{<token list>}</i>
<u><u>\quark_if_nil_p:(o V)</u></u> *	<u><u>\quark_if_nil:n</u></u> <i>TF</i> <i>{<token list>}</i> <i>{<true code>}</i> <i>{<false code>}</i>
<u><u>\quark_if_nil:n</u></u> <i>TF</i> *	Tests if the <i><token list></i> contains only <code>\q_nil</code> (distinct from <i><token list></i> being empty or containing <code>\q_nil</code> plus one or more other tokens).
<u><u>\quark_if_nil:(o V)</u></u> <i>TF</i> *	
<u><u>\quark_if_no_value_p:N</u></u> *	<u><u>\quark_if_no_value_p:N</u></u> <i><token></i>
<u><u>\quark_if_no_value_p:c</u></u> *	<u><u>\quark_if_no_value:N</u></u> <i>TF</i> <i><token></i> <i>{<true code>}</i> <i>{<false code>}</i>
<u><u>\quark_if_no_value:N</u></u> <i>TF</i> *	Tests if the <i><token></i> is equal to <code>\q_no_value</code> .
<u><u>\quark_if_no_value:c</u></u> <i>TF</i> *	
<u><u>\quark_if_no_value_p:n</u></u> *	<u><u>\quark_if_no_value_p:n</u></u> <i>{<token list>}</i>
<u><u>\quark_if_no_value:n</u></u> <i>TF</i> *	<u><u>\quark_if_no_value:n</u></u> <i>TF</i> <i>{<token list>}</i> <i>{<true code>}</i> <i>{<false code>}</i>
	Tests if the <i><token list></i> contains only <code>\q_no_value</code> (distinct from <i><token list></i> being empty or containing <code>\q_no_value</code> plus one or more other tokens).

4 Recursion

This module provides a uniform interface to intercepting and terminating loops as when one is doing tail recursion. The building blocks follow below and an example is shown in Section 4.1.

<u><u>\q_recursion_tail</u></u>	This quark is appended to the data structure in question and appears as a real element there. This means it gets any list separators around it.
---------------------------------	---

<code>\q_recursion_stop</code>	This quark is added <i>after</i> the data structure. Its purpose is to make it possible to terminate the recursion at any point easily.
--------------------------------	---

<code>\quark_if_recursion_tail_stop:N *</code>	<code>\quark_if_recursion_tail_stop:N <token></code>
--	--

Tests if $\langle token \rangle$ contains only the marker `\q_recursion_tail`, and if so uses `\use_none_delimit_by_q_recursion_stop:w` to terminate the recursion that this belongs to. The recursion input must include the marker tokens `\q_recursion_tail` and `\q_recursion_stop` as the last two items.

<code>\quark_if_recursion_tail_stop:n *</code>	<code>\quark_if_recursion_tail_stop:n {<token list>}</code>
<code>\quark_if_recursion_tail_stop:o *</code>	

Updated: 2011-09-06

Tests if the $\langle token list \rangle$ contains only `\q_recursion_tail`, and if so uses `\use_none_delimit_by_q_recursion_stop:w` to terminate the recursion that this belongs to. The recursion input must include the marker tokens `\q_recursion_tail` and `\q_recursion_stop` as the last two items.

<code>\quark_if_recursion_tail_stop_do:Nn *</code>	<code>\quark_if_recursion_tail_stop_do:Nn <token> {<insertion>}</code>
--	--

Tests if $\langle token \rangle$ contains only the marker `\q_recursion_tail`, and if so uses `\use_i_delimit_by_q_recursion_stop:w` to terminate the recursion that this belongs to. The recursion input must include the marker tokens `\q_recursion_tail` and `\q_recursion_stop` as the last two items. The $\langle insertion \rangle$ code is then added to the input stream after the recursion has ended.

<code>\quark_if_recursion_tail_stop_do:nn *</code>	<code>\quark_if_recursion_tail_stop_do:nn {<token list>} {<insertion>}</code>
<code>\quark_if_recursion_tail_stop_do:on *</code>	

Updated: 2011-09-06

Tests if the $\langle token list \rangle$ contains only `\q_recursion_tail`, and if so uses `\use_i_delimit_by_q_recursion_stop:w` to terminate the recursion that this belongs to. The recursion input must include the marker tokens `\q_recursion_tail` and `\q_recursion_stop` as the last two items. The $\langle insertion \rangle$ code is then added to the input stream after the recursion has ended.

<code>\quark_if_recursion_tail_break:NN *</code>	<code>\quark_if_recursion_tail_break:nN {<token list>}</code>
<code>\quark_if_recursion_tail_break:nN *</code>	<code>\<type>_map_break:</code>

New: 2018-04-10

Tests if $\langle token list \rangle$ contains only `\q_recursion_tail`, and if so terminates the recursion using `\<type>_map_break:.` The recursion end should be marked by `\prg_break_point:Nn \<type>_map_break:.`

4.1 An example of recursion with quarks

Quarks are mainly used internally in the `expl3` code to define recursion functions such as `\tl_map_inline:nn` and so on. Here is a small example to demonstrate how to

use quarks in this fashion. We shall define a command called `\my_map_dbl:nn` which takes a token list and applies an operation to every *pair* of tokens. For example, `\my_map_dbl:nn {abcd} {[--#1--#2--]~}` would produce “[-a-b-] [-c-d-]”. Using quarks to define such functions simplifies their logic and ensures robustness in many cases.

Here’s the definition of `\my_map_dbl:nn`. First of all, define the function that does the processing based on the inline function argument `#2`. Then initiate the recursion using an internal function. The token list `#1` is terminated using `\q_recursion_tail`, with delimiters according to the type of recursion (here a pair of `\q_recursion_tail`), concluding with `\q_recursion_stop`. These quarks are used to mark the end of the token list being operated upon.

```
\cs_new:Npn \my_map_dbl:nn #1#2
{
  \cs_set:Npn \__my_map_dbl_fn:nn ##1 ##2 {#2}
  \__my_map_dbl:nn #1 \q_recursion_tail \q_recursion_tail
  \q_recursion_stop
}
```

The definition of the internal recursion function follows. First check if either of the input tokens are the termination quarks. Then, if not, apply the inline function to the two arguments.

```
\cs_new:Nn \__my_map_dbl:nn
{
  \quark_if_recursion_tail_stop:n {#1}
  \quark_if_recursion_tail_stop:n {#2}
  \__my_map_dbl_fn:nn {#1} {#2}
}
```

Finally, recurse:

```
\__my_map_dbl:nn
}
```

Note that contrarily to L^AT_EX3 built-in mapping functions, this mapping function cannot be nested, since the second map would overwrite the definition of `__my_map_dbl_fn:nn`.

5 Scan marks

Scan marks are control sequences set equal to `\scan_stop:`, hence never expand in an expansion context and are (largely) invisible if they are encountered in a typesetting context.

Like quarks, they can be used as delimiters in weird functions and are often safer to use for this purpose. Since they are harmless when executed by T_EX in non-expandable contexts, they can be used to mark the end of a set of instructions. This allows to skip to that point if the end of the instructions should not be performed (see l3regex).

`\scan_new:N`

`\scan_new:N` *<scan mark>*

New: 2018-04-01

Creates a new *<scan mark>* which is set equal to `\scan_stop:`. The *<scan mark>* is defined globally, and an error message is raised if the name was already taken by another scan mark.

<hr/> \s_stop <hr/>	Used at the end of a set of instructions, as a marker that can be jumped to using \use_
New: 2018-04-01	none_delimit_by_s_stop:w.

<hr/> \use_none_delimit_by_s_stop:w ★	\use_none_delimit_by_s_stop:w <i><tokens></i> \s_stop
--	---

New: 2018-04-01

Removes the *<tokens>* and **\s_stop** from the input stream. This leads to a low-level T_EX error if **\s_stop** is absent.

Part VII

The l3tl package

Token lists

T_EX works with tokens, and L^AT_EX3 therefore provides a number of functions to deal with lists of tokens. Token lists may be present directly in the argument to a function:

```
\foo:n { a collection of \tokens }
```

or may be stored in a so-called “token list variable”, which have the suffix `tl`: a token list variable can also be used as the argument to a function, for example

```
\foo:N \l_some_tl
```

In both cases, functions are available to test and manipulate the lists of tokens, and these have the module prefix `tl`. In many cases, functions which can be applied to token list variables are paired with similar functions for application to explicit lists of tokens: the two “views” of a token list are therefore collected together here.

A token list (explicit, or stored in a variable) can be seen either as a list of “items”, or a list of “tokens”. An item is whatever `\use:n` would grab as its argument: a single non-space token or a brace group, with optional leading explicit space characters (each item is thus itself a token list). A token is either a normal `N` argument, or `␣`, `{`, or `}` (assuming normal T_EX category codes). Thus for example

```
{ Hello } ~ world
```

contains six items (Hello, w, o, r, l and d), but thirteen tokens (`{`, H, e, l, l, o, `}`, `␣`, w, o, r, l and d). Functions which act on items are often faster than their analogue acting directly on tokens.

1 Creating and initialising token list variables

<code>\tl_new:N</code>	<code>\tl_new:N <tl var></code>
<code>\tl_new:c</code>	

Creates a new `<tl var>` or raises an error if the name is already taken. The declaration is global. The `<tl var>` is initially empty.

<code>\tl_const:Nn</code>	<code>\tl_const:Nn <tl var> {<token list>}</code>
<code>\tl_const:(Nx cn cx)</code>	

Creates a new constant `<tl var>` or raises an error if the name is already taken. The value of the `<tl var>` is set globally to the `<token list>`.

<code>\tl_clear:N</code>	<code>\tl_clear:N <tl var></code>
<code>\tl_clear:c</code>	
<code>\tl_gclear:N</code>	
<code>\tl_gclear:c</code>	

Clears all entries from the `<tl var>`.

<hr/>	
<code>\tl_clear_new:N</code>	<code>\tl_clear_new:N <tl var></code>
<code>\tl_clear_new:c</code>	
<code>\tl_gclear_new:N</code>	Ensures that the $\langle tl\ var\rangle$ exists globally by applying <code>\tl_new:N</code> if necessary, then applies
<code>\tl_gclear_new:c</code>	<code>\tl_(g)clear:N</code> to leave the $\langle tl\ var\rangle$ empty.
<hr/>	
<code>\tl_set_eq:NN</code>	<code>\tl_set_eq:NN <tl var_1> <tl var_2></code>
<code>\tl_set_eq:(cN Nc cc)</code>	Sets the content of $\langle tl\ var_1\rangle$ equal to that of $\langle tl\ var_2\rangle$.
<code>\tl_gset_eq:NN</code>	
<code>\tl_gset_eq:(cN Nc cc)</code>	
<hr/>	
<code>\tl_concat:NNN</code>	<code>\tl_concat:NNN <tl var_1> <tl var_2> <tl var_3></code>
<code>\tl_concat:ccc</code>	
<code>\tl_gconcat:NNN</code>	Concatenates the content of $\langle tl\ var_2\rangle$ and $\langle tl\ var_3\rangle$ together and saves the result in
<code>\tl_gconcat:ccc</code>	$\langle tl\ var_1\rangle$. The $\langle tl\ var_2\rangle$ is placed at the left side of the new token list.
<hr/>	
New: 2012-05-18	
<hr/>	
<code>\tl_if_exist_p:N *</code>	<code>\tl_if_exist_p:N <tl var></code>
<code>\tl_if_exist_p:c *</code>	<code>\tl_if_exist:NTF <tl var> {\true code} {\false code}</code>
<code>\tl_if_exist:N\overline{TF} *</code>	
<code>\tl_if_exist:c\overline{TF} *</code>	Tests whether the $\langle tl\ var\rangle$ is currently defined. This does not check that the $\langle tl\ var\rangle$ really is a token list variable.
<hr/>	
New: 2012-03-03	
<hr/>	

2 Adding data to token list variables

<hr/>	
<code>\tl_set:Nn</code>	<code>\tl_set:Nn <tl var> {\tokens}</code>
<code>\tl_set:(NV Nv No Nf Nx cn cV cv co cf cx)</code>	
<code>\tl_gset:Nn</code>	
<code>\tl_gset:(NV Nv No Nf Nx cn cV cv co cf cx)</code>	
<hr/>	
Sets $\langle tl\ var\rangle$ to contain $\langle tokens\rangle$, removing any previous content from the variable.	
<hr/>	
<code>\tl_put_left:Nn</code>	<code>\tl_put_left:Nn <tl var> {\tokens}</code>
<code>\tl_put_left:(NV No Nx cn cV co cx)</code>	
<code>\tl_gput_left:Nn</code>	
<code>\tl_gput_left:(NV No Nx cn cV co cx)</code>	
<hr/>	
Appends $\langle tokens\rangle$ to the left side of the current content of $\langle tl\ var\rangle$.	
<hr/>	
<code>\tl_put_right:Nn</code>	<code>\tl_put_right:Nn <tl var> {\tokens}</code>
<code>\tl_put_right:(NV No Nx cn cV co cx)</code>	
<code>\tl_gput_right:Nn</code>	
<code>\tl_gput_right:(NV No Nx cn cV co cx)</code>	
<hr/>	
Appends $\langle tokens\rangle$ to the right side of the current content of $\langle tl\ var\rangle$.	

3 Modifying token list variables

```
\tl_replace_once:Nnn
\tl_replace_once:cnn
\tl_greplace_once:Nnn
\tl_greplace_once:cnn
```

Updated: 2011-08-11

```
\tl_replace_once:Nnn <tl var> {{<old tokens>}} {{<new tokens>}}
```

Replaces the first (leftmost) occurrence of *<old tokens>* in the *<tl var>* with *<new tokens>*. *<Old tokens>* cannot contain `{`, `}` or `#` (more precisely, explicit character tokens with category code 1 (begin-group) or 2 (end-group), and tokens with category code 6).

```
\tl_replace_all:Nnn
\tl_replace_all:cnn
\tl_greplace_all:Nnn
\tl_greplace_all:cnn
```

Updated: 2011-08-11

```
\tl_replace_all:Nnn <tl var> {{<old tokens>}} {{<new tokens>}}
```

Replaces all occurrences of *<old tokens>* in the *<tl var>* with *<new tokens>*. *<Old tokens>* cannot contain `{`, `}` or `#` (more precisely, explicit character tokens with category code 1 (begin-group) or 2 (end-group), and tokens with category code 6). As this function operates from left to right, the pattern *<old tokens>* may remain after the replacement (see `\tl_remove_all:Nn` for an example).

```
\tl_remove_once:Nn
\tl_remove_once:cn
\tl_gremove_once:Nn
\tl_gremove_once:cn
```

Updated: 2011-08-11

```
\tl_remove_once:Nn <tl var> {{<tokens>}}
```

Removes the first (leftmost) occurrence of *<tokens>* from the *<tl var>*. *<Tokens>* cannot contain `{`, `}` or `#` (more precisely, explicit character tokens with category code 1 (begin-group) or 2 (end-group), and tokens with category code 6).

```
\tl_remove_all:Nn
\tl_remove_all:cn
\tl_gremove_all:Nn
\tl_gremove_all:cn
```

Updated: 2011-08-11

```
\tl_remove_all:Nn <tl var> {{<tokens>}}
```

Removes all occurrences of *<tokens>* from the *<tl var>*. *<Tokens>* cannot contain `{`, `}` or `#` (more precisely, explicit character tokens with category code 1 (begin-group) or 2 (end-group), and tokens with category code 6). As this function operates from left to right, the pattern *<tokens>* may remain after the removal, for instance,

```
\tl_set:Nn \l_tmpa_tl {abbccd} \tl_remove_all:Nn \l_tmpa_tl {bc}
```

results in `\l_tmpa_tl` containing `abcd`.

4 Reassigning token list category codes

These functions allow the rescanning of tokens: re-apply T_EX's tokenization process to apply category codes different from those in force when the tokens were absorbed. Whilst this functionality is supported, it is often preferable to find alternative approaches to achieving outcomes rather than rescanning tokens (for example construction of token lists token-by-token with intervening category code changes or using `\char_generate:nn`).

<code>\tl_set_rescan:Nnn</code>	<code>\tl_set_rescan:Nnn <tl var> {<setup>} {<tokens>}</code>
<code>\tl_set_rescan:(Nno Nnx cnn cno cnx)</code>	
<code>\tl_gset_rescan:Nnn</code>	
<code>\tl_gset_rescan:(Nno Nnx cnn cno cnx)</code>	
Updated: 2015-08-11	

Sets $\langle tl\ var\rangle$ to contain $\langle tokens\rangle$, applying the category code régime specified in the $\langle setup\rangle$ before carrying out the assignment. (Category codes applied to tokens not explicitly covered by the $\langle setup\rangle$ are those in force at the point of use of `\tl_set_rescan:Nnn`.) This allows the $\langle tl\ var\rangle$ to contain material with category codes other than those that apply when $\langle tokens\rangle$ are absorbed. The $\langle setup\rangle$ is run within a group and may contain any valid input, although only changes in category codes, such as uses of `\cctab_select:N`, are relevant. See also `\tl_rescan:nn`.

T_EXhackers note: The $\langle tokens\rangle$ are first turned into a string (using `\tl_to_str:n`). If the string contains one or more characters with character code `\newlinechar` (set equal to `\endlinechar` unless that is equal to 32, before the user $\langle setup\rangle$), then it is split into lines at these characters, then read as if reading multiple lines from a file, ignoring spaces (catcode 10) at the beginning and spaces and tabs (character code 32 or 9) at the end of every line. Otherwise, spaces (and tabs) are retained at both ends of the single-line string, as if it appeared in the middle of a line read from a file.

<code>\tl_rescan:nn</code>	<code>\tl_rescan:nn {<setup>} {<tokens>}</code>
Updated: 2015-08-11	

Rescans $\langle tokens\rangle$ applying the category code régime specified in the $\langle setup\rangle$, and leaves the resulting tokens in the input stream. (Category codes applied to tokens not explicitly covered by the $\langle setup\rangle$ are those in force at the point of use of `\tl_rescan:nn`.) The $\langle setup\rangle$ is run within a group and may contain any valid input, although only changes in category codes, such as uses of `\cctab_select:N`, are relevant. See also `\tl_set_rescan:Nnn`, which is more robust than using `\tl_set:Nn` in the $\langle tokens\rangle$ argument of `\tl_rescan:nn`.

T_EXhackers note: The $\langle tokens\rangle$ are first turned into a string (using `\tl_to_str:n`). If the string contains one or more characters with character code `\newlinechar` (set equal to `\endlinechar` unless that is equal to 32, before the user $\langle setup\rangle$), then it is split into lines at these characters, then read as if reading multiple lines from a file, ignoring spaces (catcode 10) at the beginning and spaces and tabs (character code 32 or 9) at the end of every line. Otherwise, spaces (and tabs) are retained at both ends of the single-line string, as if it appeared in the middle of a line read from a file.

5 Token list conditionals

<hr/>	
<code>\tl_if_blank_p:n</code> ★	<code>\tl_if_blank_p:n {⟨token list⟩}</code>
<code>\tl_if_blank_p:(e V o)</code> ★	<code>\tl_if_blank:nTF {⟨token list⟩} {⟨true code⟩} {⟨false code⟩}</code>
<code>\tl_if_blank:nTF</code> ★	Tests if the $\langle token list \rangle$ consists only of blank spaces (<i>i.e.</i> contains no item). The test is true if $\langle token list \rangle$ is zero or more explicit space characters (explicit tokens with character code 32 and category code 10), and is false otherwise.
<code>\tl_if_blank:(e V o)TF</code> ★	
<hr/>	
Updated: 2019-09-04	

<code>\tl_if_empty_p:N</code>	★	<code>\tl_if_empty_p:N <tl var></code>
<code>\tl_if_empty_p:c</code>	★	<code>\tl_if_empty:NTF <tl var> {<true code>} {<false code>}</code>
<code>\tl_if_empty:nTF</code>	★	Tests if the <i><token list variable></i> is entirely empty (<i>i.e.</i> contains no tokens at all).
<code>\tl_if_empty:cTF</code>	★	

<code>\tl_if_empty_p:n</code>	★	<code>\tl_if_empty_p:n {<token list>}</code>
<code>\tl_if_empty_p:(V o)</code>	★	<code>\tl_if_empty:NNTF {<token list>} {<true code>} {<false code>}</code>
<code>\tl_if_empty:nTF</code>	★	Tests if the <i><token list></i> is entirely empty (<i>i.e.</i> contains no tokens at all).
<code>\tl_if_empty:(V o)TF</code>	★	

New: 2012-05-24
Updated: 2012-06-05

<code>\tl_if_eq_p:NN</code>	★	<code>\tl_if_eq_p:NN <tl var₁> <tl var₂></code>
<code>\tl_if_eq_p:(Nc cN cc)</code>	★	<code>\tl_if_eq:NNTF <tl var₁> <tl var₂> {<true code>} {<false code>}</code>
<code>\tl_if_eq:NNTF</code>	★	Compares the content of two <i><token list variables></i> and is logically true if the two contain the same list of tokens (<i>i.e.</i> identical in both the list of characters they contain and the category codes of those characters). Thus for example
<code>\tl_if_eq:(Nc cN cc)TF</code>	★	

```

\tl_set:Nn \l_tmpa_tl { abc }
\tl_set:Nx \l_tmpb_tl { \tl_to_str:n { abc } }
\tl_if_eq:NNTF \l_tmpa_tl \l_tmpb_tl { true } { false }

```

yields **false**. See also `\str_if_eq:nnTF` for a comparison that ignores category codes.

<code>\tl_if_eq:NnTF</code>		<code>\tl_if_eq:NnTF <tl var₁> {<token list₂>} {<true code>} {<false code>}</code>
<code>\tl_if_eq:cnTF</code>		Tests if the <i><token list variable₁></i> and the <i><token list₂></i> contain the same list of tokens, both in respect of character codes and category codes. This conditional is not expandable: see <code>\tl_if_eq:NNTF</code> for an expandable version when both token lists are stored in variables, or <code>\str_if_eq:nnTF</code> if category codes are not important.

New: 2020-07-14

<code>\tl_if_eq:nnTF</code>		<code>\tl_if_eq:nnTF {<token list₁>} {<token list₂>} {<true code>} {<false code>}</code>
		Tests if <i><token list₁></i> and <i><token list₂></i> contain the same list of tokens, both in respect of character codes and category codes. This conditional is not expandable: see <code>\tl_if_eq:NNTF</code> for an expandable version when token lists are stored in variables, or <code>\str_if_eq:nnTF</code> if category codes are not important.

<code>\tl_if_in:NnTF</code>		<code>\tl_if_in:NnTF <tl var> {<token list>} {<true code>} {<false code>}</code>
<code>\tl_if_in:cnTF</code>		Tests if the <i><token list></i> is found in the content of the <i><tl var></i> . The <i><token list></i> cannot contain the tokens <code>{</code> , <code>}</code> or <code>#</code> (more precisely, explicit character tokens with category code 1 (begin-group) or 2 (end-group), and tokens with category code 6).

<code>\tl_if_in:nnTF</code>		<code>\tl_if_in:nnTF {<token list₁>} {<token list₂>} {<true code>} {<false code>}</code>
<code>\tl_if_in:(Vn on no)TF</code>		Tests if <i><token list₂></i> is found inside <i><token list₁></i> . The <i><token list₂></i> cannot contain the tokens <code>{</code> , <code>}</code> or <code>#</code> (more precisely, explicit character tokens with category code 1 (begin-group) or 2 (end-group), and tokens with category code 6).

<code>\tl_if_novalue_p:n *</code>	<code>\tl_if_novalue_p:n {⟨token list⟩}</code>
<code>\tl_if_novalue:nTF *</code>	<code>\tl_if_novalue:nTF {⟨token list⟩} {⟨true code⟩} {⟨false code⟩}</code>

New: 2017-11-14

Tests if the $\langle token\ list \rangle$ is exactly equal to the special `\c_novalue_tl` marker. This function is intended to allow construction of flexible document interface structures in which missing optional arguments are detected.

<code>\tl_if_single_p:N *</code>	<code>\tl_if_single_p:N ⟨tl var⟩</code>
<code>\tl_if_single_p:c *</code>	<code>\tl_if_single:NTF ⟨tl var⟩ {⟨true code⟩} {⟨false code⟩}</code>
<code>\tl_if_single:NTF *</code>	
<code>\tl_if_single:cTF *</code>	

Updated: 2011-08-13

Tests if the content of the $\langle tl\ var \rangle$ consists of a single item, *i.e.* is a single normal token (neither an explicit space character nor a begin-group character) or a single brace group, surrounded by optional spaces on both sides. In other words, such a token list has token count 1 according to `\tl_count:N`.

<code>\tl_if_single_p:n *</code>	<code>\tl_if_single_p:n {⟨token list⟩}</code>
<code>\tl_if_single:nTF *</code>	<code>\tl_if_single:nTF {⟨token list⟩} {⟨true code⟩} {⟨false code⟩}</code>

Updated: 2011-08-13

Tests if the $\langle token\ list \rangle$ has exactly one item, *i.e.* is a single normal token (neither an explicit space character nor a begin-group character) or a single brace group, surrounded by optional spaces on both sides. In other words, such a token list has token count 1 according to `\tl_count:n`.

<code>\tl_if_single_token_p:n *</code>	<code>\tl_if_single_token_p:n {⟨token list⟩}</code>
<code>\tl_if_single_token:nTF *</code>	<code>\tl_if_single_token:nTF {⟨token list⟩} {⟨true code⟩} {⟨false code⟩}</code>

Tests if the token list consists of exactly one token, *i.e.* is either a single space character or a single “normal” token. Token groups $\{ \dots \}$ are not single tokens.

<code>\tl_case:Nn *</code>	<code>\tl_case:NnTF ⟨test token list variable⟩</code>
<code>\tl_case:cn *</code>	<code>{</code>
<code>\tl_case:NnTF *</code>	<code>⟨token list variable case₁⟩ {⟨code case₁⟩}</code>
<code>\tl_case:cnTF *</code>	<code>⟨token list variable case₂⟩ {⟨code case₂⟩}</code>
	<code>...</code>
	<code>⟨token list variable case_n⟩ {⟨code case_n⟩}</code>
	<code>}</code>
	<code>{⟨true code⟩}</code>
	<code>{⟨false code⟩}</code>

New: 2013-07-24

This function compares the $\langle test\ token\ list\ variable \rangle$ in turn with each of the $\langle token\ list\ variable\ cases \rangle$. If the two are equal (as described for `\tl_if_eq:NNTF`) then the associated $\langle code \rangle$ is left in the input stream and other cases are discarded. If any of the cases are matched, the $\langle true\ code \rangle$ is also inserted into the input stream (after the code for the appropriate case), while if none match then the $\langle false\ code \rangle$ is inserted. The function `\tl_case:Nn`, which does nothing if there is no match, is also available.

6 Mapping to token lists

All mappings are done at the current group level, *i.e.* any local assignments made by the $\langle function \rangle$ or $\langle code \rangle$ discussed below remain in effect after the loop.

<hr/> <code>\tl_map_function:NN</code> ☆ <code>\tl_map_function:cN</code> ☆ <hr/> Updated: 2012-06-29 <hr/>	<code>\tl_map_function:NN</code> $\langle tl\ var \rangle$ $\langle function \rangle$ Applies $\langle function \rangle$ to every $\langle item \rangle$ in the $\langle tl\ var \rangle$. The $\langle function \rangle$ receives one argument for each iteration. This may be a number of tokens if the $\langle item \rangle$ was stored within braces. Hence the $\langle function \rangle$ should anticipate receiving n-type arguments. See also <code>\tl_map_function:nN</code> .
<hr/> <code>\tl_map_function:nN</code> ☆ <hr/> Updated: 2012-06-29 <hr/>	<code>\tl_map_function:nN</code> $\{\langle token\ list \rangle\}$ $\langle function \rangle$ Applies $\langle function \rangle$ to every $\langle item \rangle$ in the $\langle token\ list \rangle$, The $\langle function \rangle$ receives one argument for each iteration. This may be a number of tokens if the $\langle item \rangle$ was stored within braces. Hence the $\langle function \rangle$ should anticipate receiving n-type arguments. See also <code>\tl_map_function:NN</code> .
<hr/> <code>\tl_map_inline:Nn</code> <code>\tl_map_inline:cn</code> <hr/> Updated: 2012-06-29 <hr/>	<code>\tl_map_inline:Nn</code> $\langle tl\ var \rangle$ $\{\langle inline\ function \rangle\}$ Applies the $\langle inline\ function \rangle$ to every $\langle item \rangle$ stored within the $\langle tl\ var \rangle$. The $\langle inline\ function \rangle$ should consist of code which receives the $\langle item \rangle$ as #1. See also <code>\tl_map_function:NN</code> .
<hr/> <code>\tl_map_inline:nn</code> <hr/> Updated: 2012-06-29 <hr/>	<code>\tl_map_inline:nn</code> $\{\langle token\ list \rangle\}$ $\{\langle inline\ function \rangle\}$ Applies the $\langle inline\ function \rangle$ to every $\langle item \rangle$ stored within the $\langle token\ list \rangle$. The $\langle inline\ function \rangle$ should consist of code which receives the $\langle item \rangle$ as #1. See also <code>\tl_map_function:nN</code> .
<hr/> <code>\tl_map_tokens:Nn</code> ☆ <code>\tl_map_tokens:cn</code> ☆ <code>\tl_map_tokens:nn</code> ☆ <hr/> New: 2019-09-02 <hr/>	<code>\tl_map_tokens:Nn</code> $\langle tl\ var \rangle$ $\{\langle code \rangle\}$ <code>\tl_map_tokens:nn</code> $\{\langle tokens \rangle\}$ $\{\langle code \rangle\}$ Analogue of <code>\tl_map_function:NN</code> which maps several tokens instead of a single function. The $\langle code \rangle$ receives each item in the $\langle tl\ var \rangle$ or in $\langle tokens \rangle$ as a trailing brace group. For instance, $\tl_map_tokens:Nn\ \l_my_tl\ \{\ \prg_replicate:nn\ \{ 2 \} \}$ expands to twice each item in the $\langle tl\ var \rangle$: for each item in <code>\l_my_tl</code> the function <code>\prg_replicate:nn</code> receives 2 and $\langle item \rangle$ as its two arguments. The function <code>\tl_map_inline:Nn</code> is typically faster but is not expandable.
<hr/> <code>\tl_map_variable:NNn</code> <code>\tl_map_variable:cNn</code> <hr/> Updated: 2012-06-29 <hr/>	<code>\tl_map_variable:NNn</code> $\langle tl\ var \rangle$ $\langle variable \rangle$ $\{\langle code \rangle\}$ Stores each $\langle item \rangle$ of the $\langle tl\ var \rangle$ in turn in the (token list) $\langle variable \rangle$ and applies the $\langle code \rangle$. The $\langle code \rangle$ will usually make use of the $\langle variable \rangle$, but this is not enforced. The assignments to the $\langle variable \rangle$ are local. Its value after the loop is the last $\langle item \rangle$ in the $\langle tl\ var \rangle$, or its original value if the $\langle tl\ var \rangle$ is blank. See also <code>\tl_map_inline:Nn</code> .
<hr/> <code>\tl_map_variable:nNn</code> <hr/> Updated: 2012-06-29 <hr/>	<code>\tl_map_variable:nNn</code> $\{\langle token\ list \rangle\}$ $\langle variable \rangle$ $\{\langle code \rangle\}$ Stores each $\langle item \rangle$ of the $\langle token\ list \rangle$ in turn in the (token list) $\langle variable \rangle$ and applies the $\langle code \rangle$. The $\langle code \rangle$ will usually make use of the $\langle variable \rangle$, but this is not enforced. The assignments to the $\langle variable \rangle$ are local. Its value after the loop is the last $\langle item \rangle$ in the $\langle tl\ var \rangle$, or its original value if the $\langle tl\ var \rangle$ is blank. See also <code>\tl_map_inline:nn</code> .

<hr/> <code>\tl_map_break:</code> ☆	<code>\tl_map_break:</code>
<hr/> Updated: 2012-06-29 <hr/>	Used to terminate a <code>\tl_map...</code> function before all entries in the <i>⟨token list variable⟩</i> have been processed. This normally takes place within a conditional statement, for example

```

\tl_map_inline:Nn \l_my_tl
{
  \str_if_eq:nnT { #1 } { bingo } { \tl_map_break: }
  % Do something useful
}

```

See also `\tl_map_break:n`. Use outside of a `\tl_map...` scenario leads to low level T_EX errors.

T_EXhackers note: When the mapping is broken, additional tokens may be inserted before the *⟨tokens⟩* are inserted into the input stream. This depends on the design of the mapping function.

<hr/> <code>\tl_map_break:n</code> ☆	<code>\tl_map_break:n {⟨code⟩}</code>
<hr/> Updated: 2012-06-29 <hr/>	Used to terminate a <code>\tl_map...</code> function before all entries in the <i>⟨token list variable⟩</i> have been processed, inserting the <i>⟨code⟩</i> after the mapping has ended. This normally takes place within a conditional statement, for example

```

\tl_map_inline:Nn \l_my_tl
{
  \str_if_eq:nnT { #1 } { bingo }
  { \tl_map_break:n { <code> } }
  % Do something useful
}

```

Use outside of a `\tl_map...` scenario leads to low level T_EX errors.

T_EXhackers note: When the mapping is broken, additional tokens may be inserted before the *⟨code⟩* is inserted into the input stream. This depends on the design of the mapping function.

7 Using token lists

<code>\tl_to_str:n</code>	★	<code>\tl_to_str:n {⟨token list⟩}</code>
<code>\tl_to_str:V</code>	★	

Converts the $\langle token\ list \rangle$ to a $\langle string \rangle$, leaving the resulting character tokens in the input stream. A $\langle string \rangle$ is a series of tokens with category code 12 (other) with the exception of spaces, which retain category code 10 (space). This function requires only a single expansion. Its argument *must* be braced.

TeXhackers note: This is the ε -TeX primitive `\detokenize`. Converting a $\langle token\ list \rangle$ to a $\langle string \rangle$ yields a concatenation of the string representations of every token in the $\langle token\ list \rangle$. The string representation of a control sequence is

- an escape character, whose character code is given by the internal parameter `\escapechar`, absent if the `\escapechar` is negative or greater than the largest character code;
- the control sequence name, as defined by `\cs_to_str:N`;
- a space, unless the control sequence name is a single character whose category at the time of expansion of `\tl_to_str:n` is not “letter”.

The string representation of an explicit character token is that character, doubled in the case of (explicit) macro parameter characters (normally `#`). In particular, the string representation of a token list may depend on the category codes in effect when it is evaluated, and the value of the `\escapechar`: for instance `\tl_to_str:n {\a}` normally produces the three character “backslash”, “lower-case a”, “space”, but it may also produce a single “lower-case a” if the escape character is negative and `a` is currently not a letter.

<code>\tl_to_str:N</code>	★	<code>\tl_to_str:N ⟨tl var⟩</code>
<code>\tl_to_str:c</code>	★	

Converts the content of the $\langle tl\ var \rangle$ into a series of characters with category code 12 (other) with the exception of spaces, which retain category code 10 (space). This $\langle string \rangle$ is then left in the input stream. For low-level details, see the notes given for `\tl_to_str:n`.

<code>\tl_use:N</code>	★	<code>\tl_use:N ⟨tl var⟩</code>
<code>\tl_use:c</code>	★	

Recovers the content of a $\langle tl\ var \rangle$ and places it directly in the input stream. An error is raised if the variable does not exist or if it is invalid. Note that it is possible to use a $\langle tl\ var \rangle$ directly without an accessor function.

8 Working with the content of token lists

<code>\tl_count:n</code>	★	<code>\tl_count:n {⟨tokens⟩}</code>
<code>\tl_count:(V o)</code>	★	

New: 2012-05-13

Counts the number of $\langle items \rangle$ in $\langle tokens \rangle$ and leaves this information in the input stream. Unbraced tokens count as one element as do each token group (`{...}`). This process ignores any unprotected spaces within $\langle tokens \rangle$. See also `\tl_count:N`. This function requires three expansions, giving an $\langle integer\ denotation \rangle$.

`\tl_count:N` ★
`\tl_count:c` ★

New: 2012-05-13

`\tl_count:N` $\langle tl\ var \rangle$
Counts the number of token groups in the $\langle tl\ var \rangle$ and leaves this information in the input stream. Unbraced tokens count as one element as do each token group $\{...\}$. This process ignores any unprotected spaces within the $\langle tl\ var \rangle$. See also `\tl_count:n`. This function requires three expansions, giving an *integer denotation*.

`\tl_count_tokens:n` ★

New: 2019-02-25

`\tl_count_tokens:n` $\{\langle tokens \rangle\}$
Counts the number of \TeX tokens in the $\langle tokens \rangle$ and leaves this information in the input stream. Every token, including spaces and braces, contributes one to the total; thus for instance, the token count of `a~{bc}` is 6.

`\tl_reverse:n` ★
`\tl_reverse:(V|o)` ★

Updated: 2012-01-08

`\tl_reverse:n` $\{\langle token\ list \rangle\}$
Reverses the order of the $\langle items \rangle$ in the $\langle token\ list \rangle$, so that $\langle item_1 \rangle \langle item_2 \rangle \langle item_3 \rangle \dots \langle item_n \rangle$ becomes $\langle item_n \rangle \dots \langle item_3 \rangle \langle item_2 \rangle \langle item_1 \rangle$. This process preserves unprotected space within the $\langle token\ list \rangle$. Tokens are not reversed within braced token groups, which keep their outer set of braces. In situations where performance is important, consider `\tl_reverse_items:n`. See also `\tl_reverse:N`.

\TeX hackers note: The result is returned within `\unexpanded`, which means that the token list does not expand further when appearing in an \mathbf{x} -type argument expansion.

`\tl_reverse:N`
`\tl_reverse:c`
`\tl_greverse:N`
`\tl_greverse:c`

Updated: 2012-01-08

`\tl_reverse:N` $\langle tl\ var \rangle$
Reverses the order of the $\langle items \rangle$ stored in $\langle tl\ var \rangle$, so that $\langle item_1 \rangle \langle item_2 \rangle \langle item_3 \rangle \dots \langle item_n \rangle$ becomes $\langle item_n \rangle \dots \langle item_3 \rangle \langle item_2 \rangle \langle item_1 \rangle$. This process preserves unprotected spaces within the $\langle token\ list\ variable \rangle$. Braced token groups are copied without reversing the order of tokens, but keep the outer set of braces. See also `\tl_reverse:n`, and, for improved performance, `\tl_reverse_items:n`.

`\tl_reverse_items:n` ★

New: 2012-01-08

`\tl_reverse_items:n` $\{\langle token\ list \rangle\}$
Reverses the order of the $\langle items \rangle$ stored in $\langle tl\ var \rangle$, so that $\{\langle item_1 \rangle\} \{\langle item_2 \rangle\} \{\langle item_3 \rangle\} \dots \{\langle item_n \rangle\}$ becomes $\{\langle item_n \rangle\} \dots \{\langle item_3 \rangle\} \{\langle item_2 \rangle\} \{\langle item_1 \rangle\}$. This process removes any unprotected space within the $\langle token\ list \rangle$. Braced token groups are copied without reversing the order of tokens, and keep the outer set of braces. Items which are initially not braced are copied with braces in the result. In cases where preserving spaces is important, consider the slower function `\tl_reverse:n`.

\TeX hackers note: The result is returned within `\unexpanded`, which means that the token list does not expand further when appearing in an \mathbf{x} -type argument expansion.

`\tl_trim_spaces:n` ★
`\tl_trim_spaces:o` ★

New: 2011-07-09
Updated: 2012-06-25

`\tl_trim_spaces:n` $\{\langle token\ list \rangle\}$
Removes any leading and trailing explicit space characters (explicit tokens with character code 32 and category code 10) from the $\langle token\ list \rangle$ and leaves the result in the input stream.

\TeX hackers note: The result is returned within `\unexpanded`, which means that the token list does not expand further when appearing in an \mathbf{x} -type argument expansion.

<code>\tl_trim_spaces_apply:nN</code> ★	<code>\tl_trim_spaces_apply:nN</code> $\{ \langle token\ list \rangle \}$ $\langle function \rangle$
<code>\tl_trim_spaces_apply:oN</code> ★	
New: 2018-04-12	Removes any leading and trailing explicit space characters (explicit tokens with character code 32 and category code 10) from the $\langle token\ list \rangle$ and passes the result to the $\langle function \rangle$ as an <i>n</i> -type argument.

<code>\tl_trim_spaces:N</code>	<code>\tl_trim_spaces:N</code> $\langle tl\ var \rangle$
<code>\tl_trim_spaces:c</code>	
<code>\tl_gtrim_spaces:N</code>	Removes any leading and trailing explicit space characters (explicit tokens with character code 32 and category code 10) from the content of the $\langle tl\ var \rangle$. Note that this therefore <i>resets</i> the content of the variable.
<code>\tl_gtrim_spaces:c</code>	
New: 2011-07-09	

<code>\tl_sort:Nn</code>	<code>\tl_sort:Nn</code> $\langle tl\ var \rangle$ $\{ \langle comparison\ code \rangle \}$
<code>\tl_sort:cn</code>	
<code>\tl_gsort:Nn</code>	Sorts the items in the $\langle tl\ var \rangle$ according to the $\langle comparison\ code \rangle$, and assigns the result to $\langle tl\ var \rangle$. The details of sorting comparison are described in Section 1.
<code>\tl_gsort:cn</code>	
New: 2017-02-06	

<code>\tl_sort:nN</code> ★	<code>\tl_sort:nN</code> $\{ \langle token\ list \rangle \}$ $\langle conditional \rangle$
New: 2017-02-06	Sorts the items in the $\langle token\ list \rangle$, using the $\langle conditional \rangle$ to compare items, and leaves the result in the input stream. The $\langle conditional \rangle$ should have signature <code>:nnTF</code> , and return true if the two items being compared should be left in the same order, and false if the items should be swapped. The details of sorting comparison are described in Section 1.

TeXhackers note: The result is returned within `\exp_not:n`, which means that the token list does not expand further when appearing in an *x*-type or *e*-type argument expansion.

9 The first token from a token list

Functions which deal with either only the very first item (balanced text or single normal token) in a token list, or the remaining tokens.

<hr/> <code>\tl_head:N</code>	★	<code>\tl_head:n {⟨token list⟩}</code>
<code>\tl_head:n</code>	★	
<code>\tl_head:(V v f)</code>	★	Leaves in the input stream the first <i>⟨item⟩</i> in the <i>⟨token list⟩</i> , discarding the rest of the <i>⟨token list⟩</i> . All leading explicit space characters (explicit tokens with character code 32 and category code 10) are discarded; for example
<hr/> Updated: 2012-09-09 <hr/>		

`\tl_head:n { abc }`

and

`\tl_head:n { ~ abc }`

both leave `a` in the input stream. If the “head” is a brace group, rather than a single token, the braces are removed, and so

`\tl_head:n { ~ { ~ ab } c }`

yields `▯ab`. A blank *⟨token list⟩* (see `\tl_if_blank:nTF`) results in `\tl_head:n` leaving nothing in the input stream.

TeXhackers note: The result is returned within `\exp_not:n`, which means that the token list does not expand further when appearing in an *x*-type argument expansion.

<hr/> <code>\tl_head:w</code>	★	<code>\tl_head:w ⟨token list⟩ { } \q_stop</code>
<hr/>		
Leaves in the input stream the first <i>⟨item⟩</i> in the <i>⟨token list⟩</i> , discarding the rest of the <i>⟨token list⟩</i> . All leading explicit space characters (explicit tokens with character code 32 and category code 10) are discarded. A blank <i>⟨token list⟩</i> (which consists only of space characters) results in a low-level TeX error, which may be avoided by the inclusion of an empty group in the input (as shown), without the need for an explicit test. Alternatively, <code>\tl_if_blank:nF</code> may be used to avoid using the function with a “blank” argument. This function requires only a single expansion, and thus is suitable for use within an <i>o</i> -type expansion. In general, <code>\tl_head:n</code> should be preferred if the number of expansions is not critical.		

<hr/> <code>\tl_tail:N</code>	★	<code>\tl_tail:n {⟨token list⟩}</code>
<code>\tl_tail:n</code>	★	
<code>\tl_tail:(V v f)</code>	★	Discards all leading explicit space characters (explicit tokens with character code 32 and category code 10) and the first <i>⟨item⟩</i> in the <i>⟨token list⟩</i> , and leaves the remaining tokens in the input stream. Thus for example
<hr/> Updated: 2012-09-01 <hr/>		

`\tl_tail:n { a ~ {bc} d }`

and

`\tl_tail:n { ~ a ~ {bc} d }`

both leave `▯{bc}d` in the input stream. A blank *⟨token list⟩* (see `\tl_if_blank:nTF`) results in `\tl_tail:n` leaving nothing in the input stream.

TeXhackers note: The result is returned within `\exp_not:n`, which means that the token list does not expand further when appearing in an *x*-type argument expansion.

```

\tl_if_head_eq_catcode_p:nN * \tl_if_head_eq_catcode_p:nN {\token list} \test token
\tl_if_head_eq_catcode_p:oN * \tl_if_head_eq_catcode:nNTF {\token list} \test token
\tl_if_head_eq_catcode:nNTF * {\true code} {\false code}
\tl_if_head_eq_catcode:oNTF *

```

Updated: 2012-07-09

Tests if the first $\langle token \rangle$ in the $\langle token list \rangle$ has the same category code as the $\langle test token \rangle$. In the case where the $\langle token list \rangle$ is empty, the test is always **false**.

```

\tl_if_head_eq_charcode_p:nN * \tl_if_head_eq_charcode_p:nN {\token list} \test token
\tl_if_head_eq_charcode_p:fN * \tl_if_head_eq_charcode:nNTF {\token list} \test token
\tl_if_head_eq_charcode:nNTF * {\true code} {\false code}
\tl_if_head_eq_charcode:fNTF *

```

Updated: 2012-07-09

Tests if the first $\langle token \rangle$ in the $\langle token list \rangle$ has the same character code as the $\langle test token \rangle$. In the case where the $\langle token list \rangle$ is empty, the test is always **false**.

```

\tl_if_head_eq_meaning_p:nN * \tl_if_head_eq_meaning_p:nN {\token list} \test token
\tl_if_head_eq_meaning:nNTF * \tl_if_head_eq_meaning:nNTF {\token list} \test token
\tl_if_head_eq_meaning:nNTF * {\true code} {\false code}

```

Updated: 2012-07-09

Tests if the first $\langle token \rangle$ in the $\langle token list \rangle$ has the same meaning as the $\langle test token \rangle$. In the case where $\langle token list \rangle$ is empty, the test is always **false**.

```

\tl_if_head_is_group_p:n * \tl_if_head_is_group_p:n {\token list}
\tl_if_head_is_group:nTF * \tl_if_head_is_group:nTF {\token list} {\true code} {\false code}

```

New: 2012-07-08

Tests if the first $\langle token \rangle$ in the $\langle token list \rangle$ is an explicit begin-group character (with category code 1 and any character code), in other words, if the $\langle token list \rangle$ starts with a brace group. In particular, the test is **false** if the $\langle token list \rangle$ starts with an implicit token such as `\c_group_begin_token`, or if it is empty. This function is useful to implement actions on token lists on a token by token basis.

```

\tl_if_head_is_N_type_p:n * \tl_if_head_is_N_type_p:n {\token list}
\tl_if_head_is_N_type:nTF * \tl_if_head_is_N_type:nTF {\token list} {\true code} {\false code}

```

New: 2012-07-08

Tests if the first $\langle token \rangle$ in the $\langle token list \rangle$ is a normal N-type argument. In other words, it is neither an explicit space character (explicit token with character code 32 and category code 10) nor an explicit begin-group character (with category code 1 and any character code). An empty argument yields **false**, as it does not have a “normal” first token. This function is useful to implement actions on token lists on a token by token basis.

```

\tl_if_head_is_space_p:n * \tl_if_head_is_space_p:n {\token list}
\tl_if_head_is_space:nTF * \tl_if_head_is_space:nTF {\token list} {\true code} {\false code}

```

Updated: 2012-07-08

Tests if the first $\langle token \rangle$ in the $\langle token list \rangle$ is an explicit space character (explicit token with character code 12 and category code 10). In particular, the test is **false** if the $\langle token list \rangle$ starts with an implicit token such as `\c_space_token`, or if it is empty. This function is useful to implement actions on token lists on a token by token basis.

10 Using a single item

<code>\tl_item:nn</code> *	<code>\tl_item:nn {(token list)} {(integer expression)}</code>
<code>\tl_item:Nn</code> *	Indexing items in the $\langle token list \rangle$ from 1 on the left, this function evaluates the $\langle integer expression \rangle$ and leaves the appropriate item from the $\langle token list \rangle$ in the input stream. If the $\langle integer expression \rangle$ is negative, indexing occurs from the right of the token list, starting at -1 for the right-most item. If the index is out of bounds, then the function expands to nothing.
<code>\tl_item:cn</code> *	

New: 2014-07-17

TeXhackers note: The result is returned within the `\unexpanded` primitive (`\exp_not:n`), which means that the $\langle item \rangle$ does not expand further when appearing in an `x`-type argument expansion.

<code>\tl_rand_item:N</code> *	<code>\tl_rand_item:N {tl var}</code>
<code>\tl_rand_item:c</code> *	<code>\tl_rand_item:n {(token list)}</code>
<code>\tl_rand_item:n</code> *	Selects a pseudo-random item of the $\langle token list \rangle$. If the $\langle token list \rangle$ is blank, the result is empty. This is not available in older versions of XeTeX.

New: 2016-12-06

TeXhackers note: The result is returned within the `\unexpanded` primitive (`\exp_not:n`), which means that the $\langle item \rangle$ does not expand further when appearing in an `x`-type argument expansion.

<code>\tl_range:Nnn</code> \star <code>\tl_range:nnn</code> \star	<code>\tl_range:Nnn</code> $\langle \text{tl var} \rangle$ $\{\langle \text{start index} \rangle\}$ $\{\langle \text{end index} \rangle\}$ <code>\tl_range:nnn</code> $\{\langle \text{token list} \rangle\}$ $\{\langle \text{start index} \rangle\}$ $\{\langle \text{end index} \rangle\}$
--	--

New: 2017-02-17
Updated: 2017-07-15

Leaves in the input stream the items from the $\langle \text{start index} \rangle$ to the $\langle \text{end index} \rangle$ inclusive. Spaces and braces are preserved between the items returned (but never at either end of the list). Here $\langle \text{start index} \rangle$ and $\langle \text{end index} \rangle$ should be $\langle \text{integer expressions} \rangle$. For describing in detail the functions' behavior, let m and n be the start and end index respectively. If either is 0, the result is empty. A positive index means 'start counting from the left end', and a negative index means 'from the right end'. Let l be the count of the token list.

The *actual start point* is determined as $M = m$ if $m > 0$ and as $M = l + m + 1$ if $m < 0$. Similarly the *actual end point* is $N = n$ if $n > 0$ and $N = l + n + 1$ if $n < 0$. If $M > N$, the result is empty. Otherwise it consists of all items from position M to position N inclusive; for the purpose of this rule, we can imagine that the token list extends at infinity on either side, with void items at positions s for $s \leq 0$ or $s > l$.

Spaces in between items in the actual range are preserved. Spaces at either end of the token list will be removed anyway (think to the token list being passed to `\tl_trim_spaces:n` to begin with).

Thus, with $l = 7$ as in the examples below, all of the following are equivalent and result in the whole token list

```
\tl_range:nnn { abcd~{e{}}fg } { 1 } { 7 }
\tl_range:nnn { abcd~{e{}}fg } { 1 } { 12 }
\tl_range:nnn { abcd~{e{}}fg } { -7 } { 7 }
\tl_range:nnn { abcd~{e{}}fg } { -12 } { 7 }
```

Here are some more interesting examples. The calls

```
\iow_term:x { \tl_range:nnn { abcd~{e{}}fg } { 2 } { 5 } }
\tl_range:nnn { abcd~{e{}}fg } { 2 } { -3 } }
\tl_range:nnn { abcd~{e{}}fg } { -6 } { 5 } }
\tl_range:nnn { abcd~{e{}}fg } { -6 } { -3 } }
```

are all equivalent and will print `bcd{e{}}` on the terminal; similarly

```
\iow_term:x { \tl_range:nnn { abcd~{e{}}fg } { 2 } { 5 } }
\tl_range:nnn { abcd~{e{}}fg } { 2 } { -3 } }
\tl_range:nnn { abcd~{e{}}fg } { -6 } { 5 } }
\tl_range:nnn { abcd~{e{}}fg } { -6 } { -3 } }
```

are all equivalent and will print `bcd {e{}}` on the terminal (note the space in the middle). To the contrary,

```
\tl_range:nnn { abcd~{e{}}f } { 2 } { 4 }
```

will discard the space after 'd'.

If we want to get the items from, say, the third to the last in a token list $\langle \text{tl} \rangle$, the call is `\tl_range:nnn { <tl> } { 3 } { -1 }`. Similarly, for discarding the last item, we can do `\tl_range:nnn { <tl> } { 1 } { -2 }`.

For better performance, see `\tl_range_braced:nnn` and `\tl_range_unbraced:nnn`.

T_EXhackers note: The result is returned within the `\unexpanded` primitive (`\exp_not:n`), which means that the $\langle \text{item} \rangle$ does not expand further when appearing in an `x`-type argument expansion.

11 Viewing token lists

`\tl_show:N`
`\tl_show:c`

Updated: 2015-08-01

`\tl_show:N <tl var>`

Displays the content of the `<tl var>` on the terminal.

TeXhackers note: This is similar to the TeX primitive `\show`, wrapped to a fixed number of characters per line.

`\tl_show:n`

Updated: 2015-08-07

`\tl_show:n <{token list}>`

Displays the `<token list>` on the terminal.

TeXhackers note: This is similar to the ϵ -TeX primitive `\showtokens`, wrapped to a fixed number of characters per line.

`\tl_log:N`
`\tl_log:c`

New: 2014-08-22
Updated: 2015-08-01

`\tl_log:N <tl var>`

Writes the content of the `<tl var>` in the log file. See also `\tl_show:N` which displays the result in the terminal.

`\tl_log:n`

New: 2014-08-22
Updated: 2015-08-07

`\tl_log:n <{token list}>`

Writes the `<token list>` in the log file. See also `\tl_show:n` which displays the result in the terminal.

12 Constant token lists

`\c_empty_tl`

Constant that is always empty.

`\c_novalue_tl`

New: 2017-11-14

A marker for the absence of an argument. This constant `tl` can safely be typeset (*cf.* `\q_nil`), with the result being `-NoValue-`. It is important to note that `\c_novalue_tl` is constructed such that it will *not* match the simple text input `-NoValue-`, *i.e.* that

`\tl_if_eq:NnTF \c_novalue_tl { -NoValue- }`

is logically **false**. The `\c_novalue_tl` marker is intended for use in creating document-level interfaces, where it serves as an indicator that an (optional) argument was omitted. In particular, it is distinct from a simple empty `tl`.

`\c_space_tl`

An explicit space character contained in a token list (compare this with `\c_space_token`). For use where an explicit space is required.

13 Scratch token lists

<hr/> <hr/>	
<code>\l_tmpa_tl</code>	
<code>\l_tmpb_tl</code>	
<hr/> <hr/>	

Scratch token lists for local assignment. These are never used by the kernel code, and so are safe for use with any $\text{\LaTeX}3$ -defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.

<hr/> <hr/>	
<code>\g_tmpa_tl</code>	
<code>\g_tmpb_tl</code>	
<hr/> <hr/>	

Scratch token lists for global assignment. These are never used by the kernel code, and so are safe for use with any $\text{\LaTeX}3$ -defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.

Part VIII

The l3str package: Strings

TeX associates each character with a category code: as such, there is no concept of a “string” as commonly understood in many other programming languages. However, there are places where we wish to manipulate token lists while in some sense “ignoring” category codes: this is done by treating token lists as strings in a TeX sense.

A TeX string (and thus an expl3 string) is a series of characters which have category code 12 (“other”) with the exception of space characters which have category code 10 (“space”). Thus at a technical level, a TeX string is a token list with the appropriate category codes. In this documentation, these are simply referred to as strings.

String variables are simply specialised token lists, but by convention should be named with the suffix `...str`. Such variables should contain characters with category code 12 (other), except spaces, which have category code 10 (blank space). All the functions in this module which accept a token list argument first convert it to a string using `\tl_to_str:n` for internal processing, and do not treat a token list or the corresponding string representation differently.

As a string is a subset of the more general token list, it is sometimes unclear when one should be used over the other. Use a string variable for data that isn’t primarily intended for typesetting and for which a level of protection from unwanted expansion is suitable. This data type simplifies comparison of variables since there are no concerns about expansion of their contents.

The functions `\cs_to_str:N`, `\tl_to_str:n`, `\tl_to_str:N` and `\token_to_str:N` (and variants) generate strings from the appropriate input: these are documented in l3basics, l3tl and l3token, respectively.

Most expandable functions in this module come in three flavours:

- `\str_...:N`, which expect a token list or string variable as their argument;
- `\str_...:n`, taking any token list (or string) as an argument;
- `\str_..._ignore_spaces:n`, which ignores any space encountered during the operation: these functions are typically faster than those which take care of escaping spaces appropriately.

1 Building strings

`\str_new:N``\str_new:c`

New: 2015-09-18

`\str_new:N <str var>`

Creates a new `<str var>` or raises an error if the name is already taken. The declaration is global. The `<str var>` is initially empty.

`\str_const:Nn``\str_const:(NV|Nx|cn|cV|cx)`New: 2015-09-18Updated: 2018-07-28

`\str_const:Nn <str var> {<token list>}`

Creates a new constant `<str var>` or raises an error if the name is already taken. The value of the `<str var>` is set globally to the `<token list>`, converted to a string.

<code>\str_clear:N</code>	<code>\str_clear:N <str var></code>
<code>\str_clear:c</code>	
<code>\str_gclear:N</code>	Clears the content of the $\langle str var \rangle$.
<code>\str_gclear:c</code>	
<hr/>	
New: 2015-09-18	

<code>\str_clear_new:N</code>	<code>\str_clear_new:N <str var></code>
<code>\str_clear_new:c</code>	
	Ensures that the $\langle str var \rangle$ exists globally by applying <code>\str_new:N</code> if necessary, then applies <code>\str_(g)clear:N</code> to leave the $\langle str var \rangle$ empty.
<hr/>	
New: 2015-09-18	

<code>\str_set_eq:NN</code>	<code>\str_set_eq:NN <str var₁> <str var₂></code>
<code>\str_set_eq:(cN Nc cc)</code>	
<code>\str_gset_eq:NN</code>	Sets the content of $\langle str var_1 \rangle$ equal to that of $\langle str var_2 \rangle$.
<code>\str_gset_eq:(cN Nc cc)</code>	
<hr/>	
New: 2015-09-18	

<code>\str_concat:NNN</code>	<code>\str_concat:NNN <str var₁> <str var₂> <str var₃></code>
<code>\str_concat:ccc</code>	
<code>\str_gconcat:NNN</code>	Concatenates the content of $\langle str var_2 \rangle$ and $\langle str var_3 \rangle$ together and saves the result in $\langle str var_1 \rangle$. The $\langle str var_2 \rangle$ is placed at the left side of the new string variable. The $\langle str var_2 \rangle$ and $\langle str var_3 \rangle$ must indeed be strings, as this function does not convert their contents to a string.
<code>\str_gconcat:ccc</code>	
<hr/>	
New: 2017-10-08	

2 Adding data to string variables

<code>\str_set:Nn</code>	<code>\str_set:Nn <str var> {<token list>}</code>
<code>\str_set:(NV Nx cn cV cx)</code>	
<code>\str_gset:Nn</code>	Converts the $\langle token list \rangle$ to a $\langle string \rangle$, and stores the result in $\langle str var \rangle$.
<code>\str_gset:(NV Nx cn cV cx)</code>	
<hr/>	
New: 2015-09-18	
Updated: 2018-07-28	

<code>\str_put_left:Nn</code>	<code>\str_put_left:Nn <str var> {<token list>}</code>
<code>\str_put_left:(NV Nx cn cV cx)</code>	
<code>\str_gput_left:Nn</code>	
<code>\str_gput_left:(NV Nx cn cV cx)</code>	
<hr/>	
New: 2015-09-18	
Updated: 2018-07-28	

Converts the $\langle token list \rangle$ to a $\langle string \rangle$, and prepends the result to $\langle str var \rangle$. The current contents of the $\langle str var \rangle$ are not automatically converted to a string.

<code>\str_put_right:Nn</code> <code>\str_put_right:(NV Nx cn cV cx)</code> <code>\str_gput_right:Nn</code> <code>\str_gput_right:(NV Nx cn cV cx)</code>	<code>\str_put_right:Nn <str var> {(token list)}</code>
--	---

New: 2015-09-18

Updated: 2018-07-28

Converts the $\langle token list \rangle$ to a $\langle string \rangle$, and appends the result to $\langle str var \rangle$. The current contents of the $\langle str var \rangle$ are not automatically converted to a string.

3 Modifying string variables

<code>\str_replace_once:Nnn</code> <code>\str_replace_once:cnn</code> <code>\str_greplace_once:Nnn</code> <code>\str_greplace_once:cnn</code>	<code>\str_replace_once:Nnn <str var> {(old)} {(new)}</code> <p>Converts the $\langle old \rangle$ and $\langle new \rangle$ token lists to strings, then replaces the first (leftmost) occurrence of $\langle old string \rangle$ in the $\langle str var \rangle$ with $\langle new string \rangle$.</p>
--	--

New: 2017-10-08

<code>\str_replace_all:Nnn</code> <code>\str_replace_all:cnn</code> <code>\str_greplace_all:Nnn</code> <code>\str_greplace_all:cnn</code>	<code>\str_replace_all:Nnn <str var> {(old)} {(new)}</code> <p>Converts the $\langle old \rangle$ and $\langle new \rangle$ token lists to strings, then replaces all occurrences of $\langle old string \rangle$ in the $\langle str var \rangle$ with $\langle new string \rangle$. As this function operates from left to right, the pattern $\langle old string \rangle$ may remain after the replacement (see <code>\str_remove_all:Nn</code> for an example).</p>
--	--

New: 2017-10-08

<code>\str_remove_once:Nn</code> <code>\str_remove_once:cn</code> <code>\str_gremove_once:Nn</code> <code>\str_gremove_once:cn</code>	<code>\str_remove_once:Nn <str var> {(token list)}</code> <p>Converts the $\langle token list \rangle$ to a $\langle string \rangle$ then removes the first (leftmost) occurrence of $\langle string \rangle$ from the $\langle str var \rangle$.</p>
--	--

New: 2017-10-08

<code>\str_remove_all:Nn</code> <code>\str_remove_all:cn</code> <code>\str_gremove_all:Nn</code> <code>\str_gremove_all:cn</code>	<code>\str_remove_all:Nn <str var> {(token list)}</code> <p>Converts the $\langle token list \rangle$ to a $\langle string \rangle$ then removes all occurrences of $\langle string \rangle$ from the $\langle str var \rangle$. As this function operates from left to right, the pattern $\langle string \rangle$ may remain after the removal, for instance,</p>
--	---

New: 2017-10-08

```
\str_set:Nn \l_tmpa_str {abbccd} \str_remove_all:Nn \l_tmpa_str
{bc}
```

results in `\l_tmpa_str` containing `abcd`.

4 String conditionals

<code>\str_if_exist_p:N</code>	★	<code>\str_if_exist_p:N</code>	⟨ <i>str var</i> ⟩
<code>\str_if_exist_p:c</code>	★	<code>\str_if_exist:NTF</code>	⟨ <i>str var</i> ⟩ {⟨ <i>true code</i> ⟩} {⟨ <i>false code</i> ⟩}
<code>\str_if_exist:NTF</code>	★	Tests whether the ⟨ <i>str var</i> ⟩ is currently defined. This does not check that the ⟨ <i>str var</i> ⟩ really is a string.	
<code>\str_if_exist:cTF</code>	★		

New: 2015-09-18

<code>\str_if_empty_p:N</code>	★	<code>\str_if_empty_p:N</code>	⟨ <i>str var</i> ⟩
<code>\str_if_empty_p:c</code>	★	<code>\str_if_empty:NTF</code>	⟨ <i>str var</i> ⟩ {⟨ <i>true code</i> ⟩} {⟨ <i>false code</i> ⟩}
<code>\str_if_empty:NTF</code>	★	Tests if the ⟨ <i>string variable</i> ⟩ is entirely empty (<i>i.e.</i> contains no characters at all).	
<code>\str_if_empty:cTF</code>	★		

New: 2015-09-18

<code>\str_if_eq_p:NN</code>	★	<code>\str_if_eq_p:NN</code>	⟨ <i>str var</i> ₁ ⟩ ⟨ <i>str var</i> ₂ ⟩
<code>\str_if_eq_p:(Nc cN cc)</code>	★	<code>\str_if_eq:NNTF</code>	⟨ <i>str var</i> ₁ ⟩ ⟨ <i>str var</i> ₂ ⟩ {⟨ <i>true code</i> ⟩} {⟨ <i>false code</i> ⟩}
<code>\str_if_eq:NNTF</code>	★	Compares the content of two ⟨ <i>str variables</i> ⟩ and is logically true if the two contain the same characters in the same order. See <code>\tl_if_eq:NNTF</code> to compare tokens (including their category codes) rather than characters.	
<code>\str_if_eq:(Nc cN cc)TF</code>	★		

New: 2015-09-18

<code>\str_if_eq_p:nn</code>	★	<code>\str_if_eq_p:nn</code>	{⟨ <i>tl</i> ₁ ⟩} {⟨ <i>tl</i> ₂ ⟩}
<code>\str_if_eq_p:(Vn on no nV VV vn nv ee)</code>	★	<code>\str_if_eq:nnTF</code>	{⟨ <i>tl</i> ₁ ⟩} {⟨ <i>tl</i> ₂ ⟩} {⟨ <i>true code</i> ⟩} {⟨ <i>false code</i> ⟩}
<code>\str_if_eq:nnTF</code>	★		
<code>\str_if_eq:(Vn on no nV VV vn nv ee)TF</code>	★		

Updated: 2018-06-18

Compares the two ⟨*token lists*⟩ on a character by character basis (namely after converting them to strings), and is **true** if the two ⟨*strings*⟩ contain the same characters in the same order. Thus for example

`\str_if_eq_p:no { abc } { \tl_to_str:n { abc } }`

is logically **true**. See `\tl_if_eq:nnTF` to compare tokens (including their category codes) rather than characters.

<code>\str_if_in:NnTF</code>	<code>\str_if_in:NnTF</code>	⟨ <i>str var</i> ⟩ {⟨ <i>token list</i> ⟩} {⟨ <i>true code</i> ⟩} {⟨ <i>false code</i> ⟩}
<code>\str_if_in:cnTF</code>		Converts the ⟨ <i>token list</i> ⟩ to a ⟨ <i>string</i> ⟩ and tests if that ⟨ <i>string</i> ⟩ is found in the content of the ⟨ <i>str var</i> ⟩.

New: 2017-10-08

<code>\str_if_in:nnTF</code>	<code>\str_if_in:nnTF</code>	⟨ <i>tl</i> ₁ ⟩ {⟨ <i>tl</i> ₂ ⟩} {⟨ <i>true code</i> ⟩} {⟨ <i>false code</i> ⟩}
		Converts both ⟨ <i>token lists</i> ⟩ to ⟨ <i>strings</i> ⟩ and tests whether ⟨ <i>string</i> ₂ ⟩ is found inside ⟨ <i>string</i> ₁ ⟩.

New: 2017-10-08

<code>\str_case:nn</code>	★	<code>\str_case:nnTF {⟨test string⟩}</code>
<code>\str_case:(Vn on nV nv)</code>	★	{
<code>\str_case:nnTF</code>	★	{⟨string case ₁ ⟩} {⟨code case ₁ ⟩}
<code>\str_case:(Vn on nV nv)TF</code>	★	{⟨string case ₂ ⟩} {⟨code case ₂ ⟩}
		...
		{⟨string case _n ⟩} {⟨code case _n ⟩}
		}
		{⟨true code⟩}
		{⟨false code⟩}

New: 2013-07-24
Updated: 2015-02-28

Compares the *⟨test string⟩* in turn with each of the *⟨string cases⟩* (all token lists are converted to strings). If the two are equal (as described for `\str_if_eq:nnTF`) then the associated *⟨code⟩* is left in the input stream and other cases are discarded. If any of the cases are matched, the *⟨true code⟩* is also inserted into the input stream (after the code for the appropriate case), while if none match then the *⟨false code⟩* is inserted. The function `\str_case:nn`, which does nothing if there is no match, is also available.

<code>\str_case_e:nn</code>	★	<code>\str_case_e:nnTF {⟨test string⟩}</code>
<code>\str_case_e:nnTF</code>	★	{
		{⟨string case ₁ ⟩} {⟨code case ₁ ⟩}
		{⟨string case ₂ ⟩} {⟨code case ₂ ⟩}
		...
		{⟨string case _n ⟩} {⟨code case _n ⟩}
		}
		{⟨true code⟩}
		{⟨false code⟩}

New: 2018-06-19

Compares the full expansion of the *⟨test string⟩* in turn with the full expansion of the *⟨string cases⟩* (all token lists are converted to strings). If the two full expansions are equal (as described for `\str_if_eq:nnTF`) then the associated *⟨code⟩* is left in the input stream and other cases are discarded. If any of the cases are matched, the *⟨true code⟩* is also inserted into the input stream (after the code for the appropriate case), while if none match then the *⟨false code⟩* is inserted. The function `\str_case_e:nn`, which does nothing if there is no match, is also available. The *⟨test string⟩* is expanded in each comparison, and must always yield the same result: for example, random numbers must not be used within this string.

5 Mapping to strings

All mappings are done at the current group level, *i.e.* any local assignments made by the *⟨function⟩* or *⟨code⟩* discussed below remain in effect after the loop.

<code>\str_map_function:NN</code>	☆	<code>\str_map_function:NN ⟨str var⟩ ⟨function⟩</code>
<code>\str_map_function:cN</code>	☆	Applies <i>⟨function⟩</i> to every <i>⟨character⟩</i> in the <i>⟨str var⟩</i> including spaces. See also <code>\str_map_function:nN</code> .
<code>\str_map_function:nN</code>	☆	<code>\str_map_function:nN {⟨token list⟩} ⟨function⟩</code>
		Converts the <i>⟨token list⟩</i> to a <i>⟨string⟩</i> then applies <i>⟨function⟩</i> to every <i>⟨character⟩</i> in the <i>⟨string⟩</i> including spaces. See also <code>\str_map_function:NN</code> .

New: 2017-11-14

<hr/> <code>\str_map_inline:Nn</code> <code>\str_map_inline:cn</code> <hr/> New: 2017-11-14	<code>\str_map_inline:Nn <str var> {<inline function>}</code> Applies the <i><inline function></i> to every <i><character></i> in the <i><str var></i> including spaces. The <i><inline function></i> should consist of code which receives the <i><character></i> as #1. See also <code>\str_map_function:NN</code> .
<hr/> <code>\str_map_inline:nn</code> <hr/> New: 2017-11-14	<code>\str_map_inline:nn {<token list>} {<inline function>}</code> Converts the <i><token list></i> to a <i><string></i> then applies the <i><inline function></i> to every <i><character></i> in the <i><string></i> including spaces. The <i><inline function></i> should consist of code which receives the <i><character></i> as #1. See also <code>\str_map_function:NN</code> .
<hr/> <code>\str_map_variable:NNn</code> <code>\str_map_variable:cNn</code> <hr/> New: 2017-11-14	<code>\str_map_variable:NNn <str var> <variable> {<code>}</code> Stores each <i><character></i> of the <i><string></i> (including spaces) in turn in the (string or token list) <i><variable></i> and applies the <i><code></i> . The <i><code></i> will usually make use of the <i><variable></i> , but this is not enforced. The assignments to the <i><variable></i> are local. Its value after the loop is the last <i><character></i> in the <i><string></i> , or its original value if the <i><string></i> is empty. See also <code>\str_map_inline:Nn</code> .
<hr/> <code>\str_map_variable:nNn</code> <hr/> New: 2017-11-14	<code>\str_map_variable:nNn {<token list>} <variable> {<code>}</code> Converts the <i><token list></i> to a <i><string></i> then stores each <i><character></i> in the <i><string></i> (including spaces) in turn in the (string or token list) <i><variable></i> and applies the <i><code></i> . The <i><code></i> will usually make use of the <i><variable></i> , but this is not enforced. The assignments to the <i><variable></i> are local. Its value after the loop is the last <i><character></i> in the <i><string></i> , or its original value if the <i><string></i> is empty. See also <code>\str_map_inline:Nn</code> .
<hr/> <code>\str_map_break: ☆</code> <hr/> New: 2017-10-08	<code>\str_map_break:</code> Used to terminate a <code>\str_map...</code> function before all characters in the <i><string></i> have been processed. This normally takes place within a conditional statement, for example

```

\str_map_inline:Nn \l_my_str
{
  \str_if_eq:nnT { #1 } { bingo } { \str_map_break: }
  % Do something useful
}

```

See also `\str_map_break:n`. Use outside of a `\str_map...` scenario leads to low level \TeX errors.

\TeX hackers note: When the mapping is broken, additional tokens may be inserted before continuing with the code that follows the loop. This depends on the design of the mapping function.

`\str_map_break:n` ☆

New: 2017-10-08

`\str_map_break:n` {*<code>*}

Used to terminate a `\str_map...` function before all characters in the *<string>* have been processed, inserting the *<code>* after the mapping has ended. This normally takes place within a conditional statement, for example

```
\str_map_inline:Nn \l_my_str
{
  \str_if_eq:nnT { #1 } { bingo }
  { \str_map_break:n { <code> } }
  % Do something useful
}
```

Use outside of a `\str_map...` scenario leads to low level TeX errors.

TeXhackers note: When the mapping is broken, additional tokens may be inserted before the *<code>* is inserted into the input stream. This depends on the design of the mapping function.

6 Working with the content of strings

`\str_use:N` ★

`\str_use:c` ★

New: 2015-09-18

`\str_use:N` *<str var>*

Recovers the content of a *<str var>* and places it directly in the input stream. An error is raised if the variable does not exist or if it is invalid. Note that it is possible to use a *<str>* directly without an accessor function.

<code>\str_count:N</code>	★	<code>\str_count:n</code> { <i><token list></i> }
<code>\str_count:c</code>	★	
<code>\str_count:n</code>	★	
<code>\str_count_ignore_spaces:n</code>	★	

New: 2015-09-18

Leaves in the input stream the number of characters in the string representation of *<token list>*, as an integer denotation. The functions differ in their treatment of spaces. In the case of `\str_count:N` and `\str_count:n`, all characters including spaces are counted. The `\str_count_ignore_spaces:n` function leaves the number of non-space characters in the input stream.

`\str_count_spaces:N` ★

`\str_count_spaces:c` ★

`\str_count_spaces:n` ★

New: 2015-09-18

`\str_count_spaces:n` {*<token list>*}

Leaves in the input stream the number of space characters in the string representation of *<token list>*, as an integer denotation. Of course, this function has no `_ignore_spaces` variant.

<code>\str_head:N</code>	★	<code>\str_head:n {⟨token list⟩}</code>
<code>\str_head:c</code>	★	
<code>\str_head:n</code>	★	
<code>\str_head_ignore_spaces:n</code>	★	

New: 2015-09-18

Converts the $\langle token\ list \rangle$ into a $\langle string \rangle$. The first character in the $\langle string \rangle$ is then left in the input stream, with category code “other”. The functions differ if the first character is a space: `\str_head:N` and `\str_head:n` return a space token with category code 10 (blank space), while the `\str_head_ignore_spaces:n` function ignores this space character and leaves the first non-space character in the input stream. If the $\langle string \rangle$ is empty (or only contains spaces in the case of the `_ignore_spaces` function), then nothing is left on the input stream.

<code>\str_tail:N</code>	★	<code>\str_tail:n {⟨token list⟩}</code>
<code>\str_tail:c</code>	★	
<code>\str_tail:n</code>	★	
<code>\str_tail_ignore_spaces:n</code>	★	

New: 2015-09-18

Converts the $\langle token\ list \rangle$ to a $\langle string \rangle$, removes the first character, and leaves the remaining characters (if any) in the input stream, with category codes 12 and 10 (for spaces). The functions differ in the case where the first character is a space: `\str_tail:N` and `\str_tail:n` only trim that space, while `\str_tail_ignore_spaces:n` removes the first non-space character and any space before it. If the $\langle token\ list \rangle$ is empty (or blank in the case of the `_ignore_spaces` variant), then nothing is left on the input stream.

<code>\str_item:Nn</code>	★	<code>\str_item:nn {⟨token list⟩} {⟨integer expression⟩}</code>
<code>\str_item:nn</code>	★	
<code>\str_item_ignore_spaces:nn</code>	★	

New: 2015-09-18

Converts the $\langle token\ list \rangle$ to a $\langle string \rangle$, and leaves in the input stream the character in position $\langle integer\ expression \rangle$ of the $\langle string \rangle$, starting at 1 for the first (left-most) character. In the case of `\str_item:Nn` and `\str_item:nn`, all characters including spaces are taken into account. The `\str_item_ignore_spaces:nn` function skips spaces when counting characters. If the $\langle integer\ expression \rangle$ is negative, characters are counted from the end of the $\langle string \rangle$. Hence, -1 is the right-most character, *etc.*

<code>\str_range:Nnn</code>	<code>*</code>	<code>\str_range:nnn {⟨token list⟩} {⟨start index⟩} {⟨end index⟩}</code>
<code>\str_range:cnn</code>	<code>*</code>	
<code>\str_range:nnn</code>	<code>*</code>	
<code>\str_range_ignore_spaces:nnn</code>	<code>*</code>	

New: 2015-09-18

Converts the $\langle token\ list \rangle$ to a $\langle string \rangle$, and leaves in the input stream the characters from the $\langle start\ index \rangle$ to the $\langle end\ index \rangle$ inclusive. Spaces are preserved and counted as items (contrast this with `\tl_range:nnn` where spaces are not counted as items and are possibly discarded from the output).

Here $\langle start\ index \rangle$ and $\langle end\ index \rangle$ should be integer denotations. For describing in detail the functions' behavior, let m and n be the start and end index respectively. If either is 0, the result is empty. A positive index means 'start counting from the left end', a negative index means 'start counting from the right end'. Let l be the count of the token list.

The *actual start point* is determined as $M = m$ if $m > 0$ and as $M = l + m + 1$ if $m < 0$. Similarly the *actual end point* is $N = n$ if $n > 0$ and $N = l + n + 1$ if $n < 0$. If $M > N$, the result is empty. Otherwise it consists of all items from position M to position N inclusive; for the purpose of this rule, we can imagine that the token list extends at infinity on either side, with void items at positions s for $s \leq 0$ or $s > l$. For instance,

```
\iow_term:x { \str_range:nnn { abcdef } { 2 } { 5 } }
\iow_term:x { \str_range:nnn { abcdef } { -4 } { -1 } }
\iow_term:x { \str_range:nnn { abcdef } { -2 } { -1 } }
\iow_term:x { \str_range:nnn { abcdef } { 0 } { -1 } }
```

prints bcde, cdef, ef, and an empty line to the terminal. The $\langle start\ index \rangle$ must always be smaller than or equal to the $\langle end\ index \rangle$: if this is not the case then no output is generated. Thus

```
\iow_term:x { \str_range:nnn { abcdef } { 5 } { 2 } }
\iow_term:x { \str_range:nnn { abcdef } { -1 } { -4 } }
```

both yield empty strings.

The behavior of `\str_range_ignore_spaces:nnn` is similar, but spaces are removed before starting the job. The input

```
\iow_term:x { \str_range:nnn { abcdefg } { 2 } { 5 } }
\iow_term:x { \str_range:nnn { abcdefg } { 2 } { -3 } }
\iow_term:x { \str_range:nnn { abcdefg } { -6 } { 5 } }
\iow_term:x { \str_range:nnn { abcdefg } { -6 } { -3 } }

\iow_term:x { \str_range:nnn { abc~efg } { 2 } { 5 } }
\iow_term:x { \str_range:nnn { abc~efg } { 2 } { -3 } }
\iow_term:x { \str_range:nnn { abc~efg } { -6 } { 5 } }
\iow_term:x { \str_range:nnn { abc~efg } { -6 } { -3 } }

\iow_term:x { \str_range_ignore_spaces:nnn { abcdefg } { 2 } { 5 } }
\iow_term:x { \str_range_ignore_spaces:nnn { abcdefg } { 2 } { -3 } }
\iow_term:x { \str_range_ignore_spaces:nnn { abcdefg } { -6 } { 5 } }
\iow_term:x { \str_range_ignore_spaces:nnn { abcdefg } { -6 } { -3 } }
```

```

\iow_term:x { \str_range_ignore_spaces:nnn { abcd~efg } { 2 } { 5 } }
\iow_term:x { \str_range_ignore_spaces:nnn { abcd~efg } { 2 } { -3 } }
\iow_term:x { \str_range_ignore_spaces:nnn { abcd~efg } { -6 } { 5 } }
\iow_term:x { \str_range_ignore_spaces:nnn { abcd~efg } { -6 } { -3 } }

```

will print four instances of `bcde`, four instances of `bc e` and eight instances of `bcde`.

7 String manipulation

```

\str_lowercase:n * \str_lowercase:n {<tokens>}
\str_lowercase:f * \str_uppercase:n {<tokens>}
\str_uppercase:n *
\str_uppercase:f *

```

New: 2019-11-26

Converts the input `<tokens>` to their string representation, as described for `\tl_to_str:n`, and then to the lower or upper case representation using a one-to-one mapping as described by the Unicode Consortium file `UnicodeData.txt`.

These functions are intended for case changing programmatic data in places where upper/lower case distinctions are meaningful. One example would be automatically generating a function name from user input where some case changing is needed. In this situation the input is programmatic, not textual, case does have meaning and a language-independent one-to-one mapping is appropriate. For example

```

\cs_new_protected:Npn \myfunc:nn #1#2
{
  \cs_set_protected:cpn
  {
    user
    \str_uppercase:f { \tl_head:n {#1} }
    \str_lowercase:f { \tl_tail:n {#1} }
  }
  { #2 }
}

```

would be used to generate a function with an auto-generated name consisting of the upper case equivalent of the supplied name followed by the lower case equivalent of the rest of the input.

These functions should *not* be used for

- Caseless comparisons: use `\str_foldcase:n` for this situation (case folding is distinct from lower casing).
- Case changing text for typesetting: see the `\text_lowercase:n(n)`, `\text_uppercase:n(n)` and `\text_titlecase:n(n)` functions which correctly deal with context-dependence and other factors appropriate to text case changing.

T_EXhackers note: As with all `expl3` functions, the input supported by `\str_foldcase:n` is *engine-native* characters which are or interoperate with UTF-8. As such, when used with `pdfTEX` *only* the Latin alphabet characters A–Z are case-folded (*i.e.* the ASCII range which coincides with UTF-8). Full UTF-8 support is available with both `XYTEX` and `LuaTEX`.

<code>\str_foldcase:n</code> *	<code>\str_foldcase:n {<tokens>}</code>
--------------------------------	---

<code>\str_foldcase:V</code> *

New: 2019-11-26

Converts the input $\langle tokens \rangle$ to their string representation, as described for `\tl_to_str:n`, and then folds the case of the resulting $\langle string \rangle$ to remove case information. The result of this process is left in the input stream.

String folding is a process used for material such as identifiers rather than for “text”. The folding provided by `\str_foldcase:n` follows the mappings provided by the [Unicode Consortium](#), who [state](#):

Case folding is primarily used for caseless comparison of text, such as identifiers in a computer program, rather than actual text transformation. Case folding in Unicode is based on the lowercase mapping, but includes additional changes to the source text to help make it language-insensitive and consistent. As a result, case-folded text should be used solely for internal processing and generally should not be stored or displayed to the end user.

The folding approach implemented by `\str_foldcase:n` follows the “full” scheme defined by the Unicode Consortium (*e.g.* SSfolds to SS). As case-folding is a language-insensitive process, there is no special treatment of Turkic input (*i.e.* I always folds to i and not to ı).

TeXhackers note: As with all `expl3` functions, the input supported by `\str_foldcase:n` is *engine-native* characters which are or interoperate with UTF-8. As such, when used with `pdfTeX` *only* the Latin alphabet characters A–Z are case-folded (*i.e.* the ASCII range which coincides with UTF-8). Full UTF-8 support is available with both `XYTeX` and `LuaTeX`, subject only to the fact that `XYTeX` in particular has issues with characters of code above hexadecimal 0xFFFF when interacting with `\tl_to_str:n`.

8 Viewing strings

<code>\str_show:N</code>	<code>\str_show:N <str var></code>
--------------------------	--

<code>\str_show:c</code>

<code>\str_show:n</code>

Displays the content of the $\langle str var \rangle$ on the terminal.

New: 2015-09-18

<code>\str_log:N</code>	<code>\str_log:N <str var></code>
-------------------------	---

<code>\str_log:c</code>

<code>\str_log:n</code>

Writes the content of the $\langle str var \rangle$ in the log file.

New: 2019-02-15

9 Constant token lists

`\c_ampersand_str`
`\c_atsign_str`
`\c_backslash_str`
`\c_left_brace_str`
`\c_right_brace_str`
`\c_circumflex_str`
`\c_colon_str`
`\c_dollar_str`
`\c_hash_str`
`\c_percent_str`
`\c_tilde_str`
`\c_underscore_str`
`\c_zero_str`

Constant strings, containing a single character token, with category code 12.

New: 2015-09-19

Updated: 2020-12-22

10 Scratch strings

`\l_tmpa_str`
`\l_tmpb_str`

Scratch strings for local assignment. These are never used by the kernel code, and so are safe for use with any $\text{\LaTeX}3$ -defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.

`\g_tmpa_str`
`\g_tmpb_str`

Scratch strings for global assignment. These are never used by the kernel code, and so are safe for use with any $\text{\LaTeX}3$ -defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.

Part IX

The `l3str-convert` package: string encoding conversions

1 Encoding and escaping schemes

Traditionally, string encodings only specify how strings of characters should be stored as bytes. However, the resulting lists of bytes are often to be used in contexts where only a restricted subset of bytes are permitted (*e.g.*, PDF string objects, URLs). Hence, storing a string of characters is done in two steps.

- The code points (“character codes”) are expressed as bytes following a given “encoding”. This can be UTF-16, ISO 8859-1, *etc.* See Table 1 for a list of supported encodings.⁶
- Bytes are translated to \TeX tokens through a given “escaping”. Those are defined for the most part by the pdf file format. See Table 2 for a list of escaping methods supported.⁶

2 Conversion functions

`\str_set_convert:Nnnn`
`\str_gset_convert:Nnnn`

`\str_set_convert:Nnnn <str var> {<string>} {<name 1>} {<name 2>}`

This function converts the $\langle string \rangle$ from the encoding given by $\langle name 1 \rangle$ to the encoding given by $\langle name 2 \rangle$, and stores the result in the $\langle str var \rangle$. Each $\langle name \rangle$ can have the form $\langle encoding \rangle$ or $\langle encoding \rangle / \langle escaping \rangle$, where the possible values of $\langle encoding \rangle$ and $\langle escaping \rangle$ are given in Tables 1 and 2, respectively. The default escaping is to input and output bytes directly. The special case of an empty $\langle name \rangle$ indicates the use of “native” strings, 8-bit for pdf \TeX , and Unicode strings for the other two engines.

For example,

`\str_set_convert:Nnnn \l_foo_str { Hello! } { } { utf16/hex }`

results in the variable `\l_foo_str` holding the string `FEFF00480065006C006C006F0021`. This is obtained by converting each character in the (native) string `Hello!` to the UTF-16 encoding, and expressing each byte as a pair of hexadecimal digits. Note the presence of a (big-endian) byte order mark “FEFF”, which can be avoided by specifying the encoding `utf16be/hex`.

An error is raised if the $\langle string \rangle$ is not valid according to the $\langle escaping 1 \rangle$ and $\langle encoding 1 \rangle$, or if it cannot be reencoded in the $\langle encoding 2 \rangle$ and $\langle escaping 2 \rangle$ (for instance, if a character does not exist in the $\langle encoding 2 \rangle$). Erroneous input is replaced by the Unicode replacement character “FFFD”, and characters which cannot be reencoded are replaced by either the replacement character “FFFD if it exists in the $\langle encoding 2 \rangle$, or an encoding-specific replacement character, or the question mark character.

⁶Encodings and escapings will be added as they are requested.

Table 1: Supported encodings. Non-alphanumeric characters are ignored, and capital letters are lower-cased before searching for the encoding in this list.

<i>⟨Encoding⟩</i>	description
<code>utf8</code>	UTF-8
<code>utf16</code>	UTF-16, with byte-order mark
<code>utf16be</code>	UTF-16, big-endian
<code>utf16le</code>	UTF-16, little-endian
<code>utf32</code>	UTF-32, with byte-order mark
<code>utf32be</code>	UTF-32, big-endian
<code>utf32le</code>	UTF-32, little-endian
<code>iso88591, latin1</code>	ISO 8859-1
<code>iso88592, latin2</code>	ISO 8859-2
<code>iso88593, latin3</code>	ISO 8859-3
<code>iso88594, latin4</code>	ISO 8859-4
<code>iso88595</code>	ISO 8859-5
<code>iso88596</code>	ISO 8859-6
<code>iso88597</code>	ISO 8859-7
<code>iso88598</code>	ISO 8859-8
<code>iso88599, latin5</code>	ISO 8859-9
<code>iso885910, latin6</code>	ISO 8859-10
<code>iso885911</code>	ISO 8859-11
<code>iso885913, latin7</code>	ISO 8859-13
<code>iso885914, latin8</code>	ISO 8859-14
<code>iso885915, latin9</code>	ISO 8859-15
<code>iso885916, latin10</code>	ISO 8859-16
<code>clist</code>	comma-list of integers
<i>⟨empty⟩</i>	native (Unicode) string

Table 2: Supported escapings. Non-alphanumeric characters are ignored, and capital letters are lower-cased before searching for the escaping in this list.

<i>⟨Escaping⟩</i>	description
<code>bytes</code> , or <code>empty</code>	arbitrary bytes
<code>hex</code> , <code>hexadecimal</code>	byte = two hexadecimal digits
<code>name</code>	see <code>\pdfescapename</code>
<code>string</code>	see <code>\pdfescapestring</code>
<code>url</code>	encoding used in URLs

<code>\str_set_convert:NnnnTF</code>	<code>\str_set_convert:NnnnTF <str var> {<string>} {<name 1>} {<name 2>} {<true code>}</code>
<code>\str_gset_convert:NnnnTF</code>	<code>{<false code>}</code>

As `\str_set_convert:Nnnn`, converts the $\langle string \rangle$ from the encoding given by $\langle name 1 \rangle$ to the encoding given by $\langle name 2 \rangle$, and assigns the result to $\langle str var \rangle$. Contrarily to `\str_set_convert:Nnnn`, the conditional variant does not raise errors in case the $\langle string \rangle$ is not valid according to the $\langle name 1 \rangle$ encoding, or cannot be expressed in the $\langle name 2 \rangle$ encoding. Instead, the $\langle false code \rangle$ is performed.

3 Conversion by expansion (for PDF contexts)

A small number of expandable functions are provided for use in PDF string/name contexts. These *assume UTF-8* and *no escaping* in the input.

<code>\str_convert_pdfname:n *</code>	<code>\str_convert_pdfname:n <string></code>
---------------------------------------	--

As `\str_set_convert:Nnnn`, converts the $\langle string \rangle$ on a byte-by-byte basis with non-ASCII codepoints escaped using hashes.

4 Possibilities, and things to do

Encoding/escaping-related tasks.

- In X_YTeX/LuaTeX, would it be better to use the $\hat{\sim}\hat{\sim}\hat{\sim}\hat{\sim}\dots$ approach to build a string from a given list of character codes? Namely, within a group, assign 0–9a–f and all characters we want to category “other”, then assign $\hat{\sim}$ the category superscript, and use `\scantokens`.
- Change `\str_set_convert:Nnnn` to expand its last two arguments.
- Describe the internal format in the code comments. Refuse code points in [“D800,”DFFF] in the internal representation?
- Add documentation about each encoding and escaping method, and add examples.
- The `hex` unescaping should raise an error for odd-token count strings.
- Decide what bytes should be escaped in the `url` escaping. Perhaps the characters `! ’ () * - . / 0 1 2 3 4 5 6 7 8 9 _` are safe, and all other characters should be escaped?
- Automate generation of 8-bit mapping files.
- Change the framework for 8-bit encodings: for decoding from 8-bit to Unicode, use 256 integer registers; for encoding, use a tree-box.
- More encodings (see Heiko’s *stringenc*). CESU?
- More escapings: ASCII85, shell escapes, lua escapes, *etc.*?

Part X

The l3seq package

Sequences and stacks

L^AT_EX3 implements a “sequence” data type, which contain an ordered list of entries which may contain any *balanced text*. It is possible to map functions to sequences such that the function is applied to every item in the sequence.

Sequences are also used to implement stack functions in L^AT_EX3. This is achieved using a number of dedicated stack functions.

1 Creating and initialising sequences

<code>\seq_new:N</code>	<code>\seq_new:N <sequence></code>
<code>\seq_new:c</code>	

Creates a new *<sequence>* or raises an error if the name is already taken. The declaration is global. The *<sequence>* initially contains no items.

<code>\seq_clear:N</code>	<code>\seq_clear:N <sequence></code>
<code>\seq_clear:c</code>	
<code>\seq_gclear:N</code>	Clears all items from the <i><sequence></i> .
<code>\seq_gclear:c</code>	

<code>\seq_clear_new:N</code>	<code>\seq_clear_new:N <sequence></code>
<code>\seq_clear_new:c</code>	
<code>\seq_gclear_new:N</code>	Ensures that the <i><sequence></i> exists globally by applying <code>\seq_new:N</code> if necessary, then applies <code>\seq_(g)clear:N</code> to leave the <i><sequence></i> empty.
<code>\seq_gclear_new:c</code>	

<code>\seq_set_eq:NN</code>	<code>\seq_set_eq:NN <sequence₁> <sequence₂></code>
<code>\seq_set_eq:(cN Nc cc)</code>	
<code>\seq_gset_eq:NN</code>	Sets the content of <i><sequence₁></i> equal to that of <i><sequence₂></i> .
<code>\seq_gset_eq:(cN Nc cc)</code>	

<code>\seq_set_from_clist:NN</code>	<code>\seq_set_from_clist:NN <sequence> <comma-list></code>
<code>\seq_set_from_clist:(cN Nc cc)</code>	
<code>\seq_set_from_clist:Nn</code>	
<code>\seq_set_from_clist:cn</code>	
<code>\seq_gset_from_clist:NN</code>	
<code>\seq_gset_from_clist:(cN Nc cc)</code>	
<code>\seq_gset_from_clist:Nn</code>	
<code>\seq_gset_from_clist:cn</code>	

New: 2014-07-17

Converts the data in the *<comma list>* into a *<sequence>*: the original *<comma list>* is unchanged.

`\seq_const_from_clist:Nn`
`\seq_const_from_clist:cn`

New: 2017-11-28

`\seq_const_from_clist:Nn` $\langle seq\ var \rangle$ $\{\langle comma-list \rangle\}$

Creates a new constant $\langle seq\ var \rangle$ or raises an error if the name is already taken. The $\langle seq\ var \rangle$ is set globally to contain the items in the $\langle comma\ list \rangle$.

`\seq_set_split:Nnn`
`\seq_set_split:NnV`
`\seq_gset_split:Nnn`
`\seq_gset_split:NnV`

New: 2011-08-15
Updated: 2012-07-02

`\seq_set_split:Nnn` $\langle sequence \rangle$ $\{\langle delimiter \rangle\}$ $\{\langle token\ list \rangle\}$

Splits the $\langle token\ list \rangle$ into $\langle items \rangle$ separated by $\langle delimiter \rangle$, and assigns the result to the $\langle sequence \rangle$. Spaces on both sides of each $\langle item \rangle$ are ignored, then one set of outer braces is removed (if any); this space trimming behaviour is identical to that of `\l3clist` functions. Empty $\langle items \rangle$ are preserved by `\seq_set_split:Nnn`, and can be removed afterwards using `\seq_remove_all:Nn` $\langle sequence \rangle$ $\{\}$. The $\langle delimiter \rangle$ may not contain `{`, `}` or `#` (assuming \TeX 's normal category code régime). If the $\langle delimiter \rangle$ is empty, the $\langle token\ list \rangle$ is split into $\langle items \rangle$ as a $\langle token\ list \rangle$.

`\seq_concat:NNN`
`\seq_concat:ccc`
`\seq_gconcat:NNN`
`\seq_gconcat:ccc`

`\seq_concat:NNN` $\langle sequence_1 \rangle$ $\langle sequence_2 \rangle$ $\langle sequence_3 \rangle$

Concatenates the content of $\langle sequence_2 \rangle$ and $\langle sequence_3 \rangle$ together and saves the result in $\langle sequence_1 \rangle$. The items in $\langle sequence_2 \rangle$ are placed at the left side of the new sequence.

`\seq_if_exist_p:N *`
`\seq_if_exist_p:c *`
`\seq_if_exist:NTF *`
`\seq_if_exist:cTF *`

New: 2012-03-03

`\seq_if_exist_p:N` $\langle sequence \rangle$

`\seq_if_exist:NTF` $\langle sequence \rangle$ $\{\langle true\ code \rangle\}$ $\{\langle false\ code \rangle\}$

Tests whether the $\langle sequence \rangle$ is currently defined. This does not check that the $\langle sequence \rangle$ really is a sequence variable.

2 Appending data to sequences

`\seq_put_left:Nn`
`\seq_put_left:(NV|Nv|No|Nx|cn|cV|cv|co|cx)`
`\seq_gput_left:Nn`
`\seq_gput_left:(NV|Nv|No|Nx|cn|cV|cv|co|cx)`

`\seq_put_left:Nn` $\langle sequence \rangle$ $\{\langle item \rangle\}$

Appends the $\langle item \rangle$ to the left of the $\langle sequence \rangle$.

`\seq_put_right:Nn`
`\seq_put_right:(NV|Nv|No|Nx|cn|cV|cv|co|cx)`
`\seq_gput_right:Nn`
`\seq_gput_right:(NV|Nv|No|Nx|cn|cV|cv|co|cx)`

`\seq_put_right:Nn` $\langle sequence \rangle$ $\{\langle item \rangle\}$

Appends the $\langle item \rangle$ to the right of the $\langle sequence \rangle$.

3 Recovering items from sequences

Items can be recovered from either the left or the right of sequences. For implementation reasons, the actions at the left of the sequence are faster than those acting on the right. These functions all assign the recovered material locally, *i.e.* setting the $\langle token\ list\ variable \rangle$ used with `\tl_set:Nn` and *never* `\tl_gset:Nn`.

<hr/> <code>\seq_get_left:NN</code> <code>\seq_get_left:cN</code> <hr/> Updated: 2012-05-14 <hr/>	<code>\seq_get_left:NN</code> $\langle sequence \rangle$ $\langle token list variable \rangle$ Stores the left-most item from a $\langle sequence \rangle$ in the $\langle token list variable \rangle$ without removing it from the $\langle sequence \rangle$. The $\langle token list variable \rangle$ is assigned locally. If $\langle sequence \rangle$ is empty the $\langle token list variable \rangle$ is set to the special marker <code>\q_no_value</code> .
<hr/> <code>\seq_get_right:NN</code> <code>\seq_get_right:cN</code> <hr/> Updated: 2012-05-19 <hr/>	<code>\seq_get_right:NN</code> $\langle sequence \rangle$ $\langle token list variable \rangle$ Stores the right-most item from a $\langle sequence \rangle$ in the $\langle token list variable \rangle$ without removing it from the $\langle sequence \rangle$. The $\langle token list variable \rangle$ is assigned locally. If $\langle sequence \rangle$ is empty the $\langle token list variable \rangle$ is set to the special marker <code>\q_no_value</code> .
<hr/> <code>\seq_pop_left:NN</code> <code>\seq_pop_left:cN</code> <hr/> Updated: 2012-05-14 <hr/>	<code>\seq_pop_left:NN</code> $\langle sequence \rangle$ $\langle token list variable \rangle$ Pops the left-most item from a $\langle sequence \rangle$ into the $\langle token list variable \rangle$, <i>i.e.</i> removes the item from the sequence and stores it in the $\langle token list variable \rangle$. Both of the variables are assigned locally. If $\langle sequence \rangle$ is empty the $\langle token list variable \rangle$ is set to the special marker <code>\q_no_value</code> .
<hr/> <code>\seq_gpop_left:NN</code> <code>\seq_gpop_left:cN</code> <hr/> Updated: 2012-05-14 <hr/>	<code>\seq_gpop_left:NN</code> $\langle sequence \rangle$ $\langle token list variable \rangle$ Pops the left-most item from a $\langle sequence \rangle$ into the $\langle token list variable \rangle$, <i>i.e.</i> removes the item from the sequence and stores it in the $\langle token list variable \rangle$. The $\langle sequence \rangle$ is modified globally, while the assignment of the $\langle token list variable \rangle$ is local. If $\langle sequence \rangle$ is empty the $\langle token list variable \rangle$ is set to the special marker <code>\q_no_value</code> .
<hr/> <code>\seq_pop_right:NN</code> <code>\seq_pop_right:cN</code> <hr/> Updated: 2012-05-19 <hr/>	<code>\seq_pop_right:NN</code> $\langle sequence \rangle$ $\langle token list variable \rangle$ Pops the right-most item from a $\langle sequence \rangle$ into the $\langle token list variable \rangle$, <i>i.e.</i> removes the item from the sequence and stores it in the $\langle token list variable \rangle$. Both of the variables are assigned locally. If $\langle sequence \rangle$ is empty the $\langle token list variable \rangle$ is set to the special marker <code>\q_no_value</code> .
<hr/> <code>\seq_gpop_right:NN</code> <code>\seq_gpop_right:cN</code> <hr/> Updated: 2012-05-19 <hr/>	<code>\seq_gpop_right:NN</code> $\langle sequence \rangle$ $\langle token list variable \rangle$ Pops the right-most item from a $\langle sequence \rangle$ into the $\langle token list variable \rangle$, <i>i.e.</i> removes the item from the sequence and stores it in the $\langle token list variable \rangle$. The $\langle sequence \rangle$ is modified globally, while the assignment of the $\langle token list variable \rangle$ is local. If $\langle sequence \rangle$ is empty the $\langle token list variable \rangle$ is set to the special marker <code>\q_no_value</code> .
<hr/> <code>\seq_item:Nn</code> ★ <code>\seq_item:cn</code> ★ <hr/> New: 2014-07-17 <hr/>	<code>\seq_item:Nn</code> $\langle sequence \rangle$ $\{ \langle integer expression \rangle \}$ Indexing items in the $\langle sequence \rangle$ from 1 at the top (left), this function evaluates the $\langle integer expression \rangle$ and leaves the appropriate item from the sequence in the input stream. If the $\langle integer expression \rangle$ is negative, indexing occurs from the bottom (right) of the sequence. If the $\langle integer expression \rangle$ is larger than the number of items in the $\langle sequence \rangle$ (as calculated by <code>\seq_count:N</code>) then the function expands to nothing.

T_EXhackers note: The result is returned within the `\unexpanded` primitive (`\exp_not:n`), which means that the $\langle item \rangle$ does not expand further when appearing in an *x*-type argument expansion.

`\seq_rand_item:N` ★
`\seq_rand_item:c` ★

New: 2016-12-06

`\seq_rand_item:N` $\langle seq\ var \rangle$

Selects a pseudo-random item of the $\langle sequence \rangle$. If the $\langle sequence \rangle$ is empty the result is empty. This is not available in older versions of Xe_{La}TeX.

TeXhackers note: The result is returned within the `\unexpanded` primitive (`\exp_not:n`), which means that the $\langle item \rangle$ does not expand further when appearing in an x-type argument expansion.

4 Recovering values from sequences with branching

The functions in this section combine tests for non-empty sequences with recovery of an item from the sequence. They offer increased readability and performance over separate testing and recovery phases.

`\seq_get_left:NNTF`
`\seq_get_left:cNTF`

New: 2012-05-14
Updated: 2012-05-19

`\seq_get_left:NNTF` $\langle sequence \rangle$ $\langle token\ list\ variable \rangle$ $\{\langle true\ code \rangle\}$ $\{\langle false\ code \rangle\}$

If the $\langle sequence \rangle$ is empty, leaves the $\langle false\ code \rangle$ in the input stream. The value of the $\langle token\ list\ variable \rangle$ is not defined in this case and should not be relied upon. If the $\langle sequence \rangle$ is non-empty, stores the left-most item from the $\langle sequence \rangle$ in the $\langle token\ list\ variable \rangle$ without removing it from the $\langle sequence \rangle$, then leaves the $\langle true\ code \rangle$ in the input stream. The $\langle token\ list\ variable \rangle$ is assigned locally.

`\seq_get_right:NNTF`
`\seq_get_right:cNTF`

New: 2012-05-19

`\seq_get_right:NNTF` $\langle sequence \rangle$ $\langle token\ list\ variable \rangle$ $\{\langle true\ code \rangle\}$ $\{\langle false\ code \rangle\}$

If the $\langle sequence \rangle$ is empty, leaves the $\langle false\ code \rangle$ in the input stream. The value of the $\langle token\ list\ variable \rangle$ is not defined in this case and should not be relied upon. If the $\langle sequence \rangle$ is non-empty, stores the right-most item from the $\langle sequence \rangle$ in the $\langle token\ list\ variable \rangle$ without removing it from the $\langle sequence \rangle$, then leaves the $\langle true\ code \rangle$ in the input stream. The $\langle token\ list\ variable \rangle$ is assigned locally.

`\seq_pop_left:NNTF`
`\seq_pop_left:cNTF`

New: 2012-05-14
Updated: 2012-05-19

`\seq_pop_left:NNTF` $\langle sequence \rangle$ $\langle token\ list\ variable \rangle$ $\{\langle true\ code \rangle\}$ $\{\langle false\ code \rangle\}$

If the $\langle sequence \rangle$ is empty, leaves the $\langle false\ code \rangle$ in the input stream. The value of the $\langle token\ list\ variable \rangle$ is not defined in this case and should not be relied upon. If the $\langle sequence \rangle$ is non-empty, pops the left-most item from the $\langle sequence \rangle$ in the $\langle token\ list\ variable \rangle$, *i.e.* removes the item from the $\langle sequence \rangle$, then leaves the $\langle true\ code \rangle$ in the input stream. Both the $\langle sequence \rangle$ and the $\langle token\ list\ variable \rangle$ are assigned locally.

`\seq_gpop_left:NNTF`
`\seq_gpop_left:cNTF`

New: 2012-05-14
Updated: 2012-05-19

`\seq_gpop_left:NNTF` $\langle sequence \rangle$ $\langle token\ list\ variable \rangle$ $\{\langle true\ code \rangle\}$ $\{\langle false\ code \rangle\}$

If the $\langle sequence \rangle$ is empty, leaves the $\langle false\ code \rangle$ in the input stream. The value of the $\langle token\ list\ variable \rangle$ is not defined in this case and should not be relied upon. If the $\langle sequence \rangle$ is non-empty, pops the left-most item from the $\langle sequence \rangle$ in the $\langle token\ list\ variable \rangle$, *i.e.* removes the item from the $\langle sequence \rangle$, then leaves the $\langle true\ code \rangle$ in the input stream. The $\langle sequence \rangle$ is modified globally, while the $\langle token\ list\ variable \rangle$ is assigned locally.

<code>\seq_pop_right:nnTF</code>	<code>\seq_pop_right:nnTF <sequence> <token list variable> {<true code>} {<false code>}</code>
<code>\seq_pop_right:cnnTF</code>	
New: 2012-05-19	

If the *<sequence>* is empty, leaves the *<false code>* in the input stream. The value of the *<token list variable>* is not defined in this case and should not be relied upon. If the *<sequence>* is non-empty, pops the right-most item from the *<sequence>* in the *<token list variable>*, *i.e.* removes the item from the *<sequence>*, then leaves the *<true code>* in the input stream. Both the *<sequence>* and the *<token list variable>* are assigned locally.

<code>\seq_gpop_right:nnTF</code>	<code>\seq_gpop_right:nnTF <sequence> <token list variable> {<true code>} {<false code>}</code>
<code>\seq_gpop_right:cnnTF</code>	
New: 2012-05-19	

If the *<sequence>* is empty, leaves the *<false code>* in the input stream. The value of the *<token list variable>* is not defined in this case and should not be relied upon. If the *<sequence>* is non-empty, pops the right-most item from the *<sequence>* in the *<token list variable>*, *i.e.* removes the item from the *<sequence>*, then leaves the *<true code>* in the input stream. The *<sequence>* is modified globally, while the *<token list variable>* is assigned locally.

5 Modifying sequences

While sequences are normally used as ordered lists, it may be necessary to modify the content. The functions here may be used to update sequences, while retaining the order of the unaffected entries.

<code>\seq_remove_duplicates:n</code>	<code>\seq_remove_duplicates:n <sequence></code>
<code>\seq_remove_duplicates:c</code>	
<code>\seq_gremove_duplicates:n</code>	Removes duplicate items from the <i><sequence></i> , leaving the left most copy of each item in the <i><sequence></i> . The <i><item></i> comparison takes place on a token basis, as for <code>\tl_if_eq:nnTF</code> .
<code>\seq_gremove_duplicates:c</code>	

TeXhackers note: This function iterates through every item in the *<sequence>* and does a comparison with the *<items>* already checked. It is therefore relatively slow with large sequences.

<code>\seq_remove_all:nn</code>	<code>\seq_remove_all:nn <sequence> {<item>}</code>
<code>\seq_remove_all:cn</code>	
<code>\seq_gremove_all:nn</code>	Removes every occurrence of <i><item></i> from the <i><sequence></i> . The <i><item></i> comparison takes place on a token basis, as for <code>\tl_if_eq:nnTF</code> .
<code>\seq_gremove_all:cn</code>	

<code>\seq_reverse:n</code>	<code>\seq_reverse:n <sequence></code>
<code>\seq_reverse:c</code>	
<code>\seq_greverse:n</code>	Reverses the order of the items stored in the <i><sequence></i> .
<code>\seq_greverse:c</code>	

New: 2014-07-18

<code>\seq_sort:nn</code>	<code>\seq_sort:nn <sequence> {<comparison code>}</code>
<code>\seq_sort:cn</code>	
<code>\seq_gsort:nn</code>	Sorts the items in the <i><sequence></i> according to the <i><comparison code></i> , and assigns the result to <i><sequence></i> . The details of sorting comparison are described in Section 1.
<code>\seq_gsort:cn</code>	

New: 2017-02-06

```
\seq_shuffle:N
\seq_shuffle:c
\seq_gshuffle:N
\seq_gshuffle:c
```

New: 2018-04-29

```
\seq_shuffle:N <seq var>
```

Sets the $\langle seq\ var \rangle$ to the result of placing the items of the $\langle seq\ var \rangle$ in a random order. Each item is (roughly) as likely to end up in any given position.

T_EXhackers note: For sequences with more than 13 items or so, only a small proportion of all possible permutations can be reached, because the random seed `\sys_rand_seed` only has 28-bits. The use of `\toks` internally means that sequences with more than 32767 or 65535 items (depending on the engine) cannot be shuffled.

6 Sequence conditionals

```
\seq_if_empty_p:N *
\seq_if_empty_p:c *
\seq_if_empty:NTF *
\seq_if_empty:cTF *
```

```
\seq_if_empty_p:N <sequence>
```

```
\seq_if_empty:NTF <sequence> {\true code} {\false code}
```

Tests if the $\langle sequence \rangle$ is empty (containing no items).

```
\seq_if_in:NnTF
```

```
\seq_if_in:NnTF <sequence> {\item} {\true code} {\false code}
```

```
\seq_if_in:(NV|Nv|No|Nx|cn|cV|cv|co|cx)TF
```

Tests if the $\langle item \rangle$ is present in the $\langle sequence \rangle$.

7 Mapping to sequences

All mappings are done at the current group level, *i.e.* any local assignments made by the $\langle function \rangle$ or $\langle code \rangle$ discussed below remain in effect after the loop.

```
\seq_map_function:NN ☆
\seq_map_function:cN ☆
```

Updated: 2012-06-29

```
\seq_map_function:NN <sequence> <function>
```

Applies $\langle function \rangle$ to every $\langle item \rangle$ stored in the $\langle sequence \rangle$. The $\langle function \rangle$ will receive one argument for each iteration. The $\langle items \rangle$ are returned from left to right. To pass further arguments to the $\langle function \rangle$, see `\seq_map_tokens:Nn`. The function `\seq_map_inline:Nn` is faster than `\seq_map_function:NN` for sequences with more than about 10 items.

```
\seq_map_inline:Nn
\seq_map_inline:cn
```

Updated: 2012-06-29

```
\seq_map_inline:Nn <sequence> {\inline function}
```

Applies $\langle inline\ function \rangle$ to every $\langle item \rangle$ stored within the $\langle sequence \rangle$. The $\langle inline\ function \rangle$ should consist of code which will receive the $\langle item \rangle$ as #1. The $\langle items \rangle$ are returned from left to right.

`\seq_map_tokens:Nn` ☆

`\seq_map_tokens:cn` ☆

New: 2019-08-30

`\seq_map_tokens:Nn` $\langle sequence \rangle$ $\{ \langle code \rangle \}$

Analogue of `\seq_map_function:NN` which maps several tokens instead of a single function. The $\langle code \rangle$ receives each item in the $\langle sequence \rangle$ as a trailing brace group. For instance,

`\seq_map_tokens:Nn \l_my_seq { \prg_replicate:nn { 2 } }`

expands to twice each item in the $\langle sequence \rangle$: for each item in `\l_my_seq` the function `\prg_replicate:nn` receives 2 and $\langle item \rangle$ as its two arguments. The function `\seq_map_inline:Nn` is typically faster but it is not expandable.

`\seq_map_variable:NNn`

`\seq_map_variable:(Ncn|cNn|ccn)`

Updated: 2012-06-29

`\seq_map_variable:NNn` $\langle sequence \rangle$ $\langle variable \rangle$ $\{ \langle code \rangle \}$

Stores each $\langle item \rangle$ of the $\langle sequence \rangle$ in turn in the (token list) $\langle variable \rangle$ and applies the $\langle code \rangle$. The $\langle code \rangle$ will usually make use of the $\langle variable \rangle$, but this is not enforced. The assignments to the $\langle variable \rangle$ are local. Its value after the loop is the last $\langle item \rangle$ in the $\langle sequence \rangle$, or its original value if the $\langle sequence \rangle$ is empty. The $\langle items \rangle$ are returned from left to right.

`\seq_map_indexed_function:NN` ☆

`\seq_map_indexed_function:NN` $\langle seq var \rangle$ $\langle function \rangle$

New: 2018-05-03

Applies $\langle function \rangle$ to every entry in the $\langle sequence variable \rangle$. The $\langle function \rangle$ should have signature `:nn`. It receives two arguments for each iteration: the $\langle index \rangle$ (namely 1 for the first entry, then 2 and so on) and the $\langle item \rangle$.

`\seq_map_indexed_inline:Nn`

New: 2018-05-03

`\seq_map_indexed_inline:Nn` $\langle seq var \rangle$ $\{ \langle inline function \rangle \}$

Applies $\langle inline function \rangle$ to every entry in the $\langle sequence variable \rangle$. The $\langle inline function \rangle$ should consist of code which receives the $\langle index \rangle$ (namely 1 for the first entry, then 2 and so on) as #1 and the $\langle item \rangle$ as #2.

\seq_map_break: ☆

Updated: 2012-06-29

\seq_map_break:

Used to terminate a `\seq_map...` function before all entries in the $\langle sequence \rangle$ have been processed. This normally takes place within a conditional statement, for example

```
\seq_map_inline:Nn \l_my_seq
{
  \str_if_eq:nnTF { #1 } { bingo }
  { \seq_map_break: }
  {
    % Do something useful
  }
}
```

Use outside of a `\seq_map...` scenario leads to low level T_EX errors.

T_EXhackers note: When the mapping is broken, additional tokens may be inserted before further items are taken from the input stream. This depends on the design of the mapping function.

\seq_map_break:n ☆

Updated: 2012-06-29

\seq_map_break:n $\{\langle code \rangle\}$

Used to terminate a `\seq_map...` function before all entries in the $\langle sequence \rangle$ have been processed, inserting the $\langle code \rangle$ after the mapping has ended. This normally takes place within a conditional statement, for example

```
\seq_map_inline:Nn \l_my_seq
{
  \str_if_eq:nnTF { #1 } { bingo }
  { \seq_map_break:n { <code> } }
  {
    % Do something useful
  }
}
```

Use outside of a `\seq_map...` scenario leads to low level T_EX errors.

T_EXhackers note: When the mapping is broken, additional tokens may be inserted before the $\langle code \rangle$ is inserted into the input stream. This depends on the design of the mapping function.

\seq_set_map:NNn
\seq_gset_map:NNn

New: 2011-12-22
Updated: 2020-07-16

\seq_set_map:NNn $\langle sequence_1 \rangle$ $\langle sequence_2 \rangle$ $\{\langle inline function \rangle\}$

Applies $\langle inline function \rangle$ to every $\langle item \rangle$ stored within the $\langle sequence_2 \rangle$. The $\langle inline function \rangle$ should consist of code which will receive the $\langle item \rangle$ as #1. The sequence resulting applying $\langle inline function \rangle$ to each $\langle item \rangle$ is assigned to $\langle sequence_1 \rangle$.

T_EXhackers note: Contrarily to other mapping functions, `\seq_map_break:` cannot be used in this function, and would lead to low-level T_EX errors.

`\seq_set_map_x:NNn`
`\seq_gset_map_x:NNn`

New: 2020-07-16

`\seq_set_map_x:NNn` $\langle sequence_1 \rangle$ $\langle sequence_2 \rangle$ $\{\langle inline function \rangle\}$

Applies $\langle inline function \rangle$ to every $\langle item \rangle$ stored within the $\langle sequence_2 \rangle$. The $\langle inline function \rangle$ should consist of code which will receive the $\langle item \rangle$ as #1. The sequence resulting from x-expanding $\langle inline function \rangle$ applied to each $\langle item \rangle$ is assigned to $\langle sequence_1 \rangle$. As such, the code in $\langle inline function \rangle$ should be expandable.

T_EXhackers note: Contrarily to other mapping functions, `\seq_map_break:` cannot be used in this function, and would lead to low-level T_EX errors.

`\seq_count:N` ★
`\seq_count:c` ★

New: 2012-07-13

`\seq_count:N` $\langle sequence \rangle$

Leaves the number of items in the $\langle sequence \rangle$ in the input stream as an $\langle integer denotation \rangle$. The total number of items in a $\langle sequence \rangle$ includes those which are empty and duplicates, *i.e.* every item in a $\langle sequence \rangle$ is unique.

8 Using the content of sequences directly

`\seq_use:Nnnn` ★
`\seq_use:cnnn` ★

New: 2013-05-26

`\seq_use:Nnnn` $\langle seq var \rangle$ $\{\langle separator between two \rangle\}$
 $\{\langle separator between more than two \rangle\}$ $\{\langle separator between final two \rangle\}$

Places the contents of the $\langle seq var \rangle$ in the input stream, with the appropriate $\langle separator \rangle$ between the items. Namely, if the sequence has more than two items, the $\langle separator between more than two \rangle$ is placed between each pair of items except the last, for which the $\langle separator between final two \rangle$ is used. If the sequence has exactly two items, then they are placed in the input stream separated by the $\langle separator between two \rangle$. If the sequence has a single item, it is placed in the input stream, and an empty sequence produces no output. An error is raised if the variable does not exist or if it is invalid.

For example,

```
\seq_set_split:Nnn \l_tmpa_seq { | } { a | b | c | {de} | f }
\seq_use:Nnnn \l_tmpa_seq { ~and~ } { ,~ } { ,~and~ }
```

inserts “a, b, c, de, and f” in the input stream. The first separator argument is not used in this case because the sequence has more than 2 items.

T_EXhackers note: The result is returned within the `\unexpanded` primitive (`\exp_not:n`), which means that the $\langle items \rangle$ do not expand further when appearing in an x-type argument expansion.

`\seq_use:Nn` ★
`\seq_use:cn` ★

New: 2013-05-26

`\seq_use:Nn` $\langle seq\ var \rangle$ $\{\langle separator \rangle\}$

Places the contents of the $\langle seq\ var \rangle$ in the input stream, with the $\langle separator \rangle$ between the items. If the sequence has a single item, it is placed in the input stream with no $\langle separator \rangle$, and an empty sequence produces no output. An error is raised if the variable does not exist or if it is invalid.

For example,

```
\seq_set_split:Nnn \l_tmpa_seq { | } { a | b | c | {de} | f }
\seq_use:Nn \l_tmpa_seq { ~and~ }
```

inserts “a and b and c and de and f” in the input stream.

TeXhackers note: The result is returned within the `\unexpanded` primitive (`\exp_not:n`), which means that the $\langle items \rangle$ do not expand further when appearing in an `x`-type argument expansion.

9 Sequences as stacks

Sequences can be used as stacks, where data is pushed to and popped from the top of the sequence. (The left of a sequence is the top, for performance reasons.) The stack functions for sequences are not intended to be mixed with the general ordered data functions detailed in the previous section: a sequence should either be used as an ordered data type or as a stack, but not in both ways.

`\seq_get:NN`
`\seq_get:cn`

Updated: 2012-05-14

`\seq_get:NN` $\langle sequence \rangle$ $\langle token\ list\ variable \rangle$

Reads the top item from a $\langle sequence \rangle$ into the $\langle token\ list\ variable \rangle$ without removing it from the $\langle sequence \rangle$. The $\langle token\ list\ variable \rangle$ is assigned locally. If $\langle sequence \rangle$ is empty the $\langle token\ list\ variable \rangle$ is set to the special marker `\q_no_value`.

`\seq_pop:NN`
`\seq_pop:cn`

Updated: 2012-05-14

`\seq_pop:NN` $\langle sequence \rangle$ $\langle token\ list\ variable \rangle$

Pops the top item from a $\langle sequence \rangle$ into the $\langle token\ list\ variable \rangle$. Both of the variables are assigned locally. If $\langle sequence \rangle$ is empty the $\langle token\ list\ variable \rangle$ is set to the special marker `\q_no_value`.

`\seq_gpop:NN`
`\seq_gpop:cn`

Updated: 2012-05-14

`\seq_gpop:NN` $\langle sequence \rangle$ $\langle token\ list\ variable \rangle$

Pops the top item from a $\langle sequence \rangle$ into the $\langle token\ list\ variable \rangle$. The $\langle sequence \rangle$ is modified globally, while the $\langle token\ list\ variable \rangle$ is assigned locally. If $\langle sequence \rangle$ is empty the $\langle token\ list\ variable \rangle$ is set to the special marker `\q_no_value`.

`\seq_get:NNTF`
`\seq_get:cNTF`

New: 2012-05-14
Updated: 2012-05-19

`\seq_get:NNTF` $\langle sequence \rangle$ $\langle token\ list\ variable \rangle$ $\{\langle true\ code \rangle\}$ $\{\langle false\ code \rangle\}$

If the $\langle sequence \rangle$ is empty, leaves the $\langle false\ code \rangle$ in the input stream. The value of the $\langle token\ list\ variable \rangle$ is not defined in this case and should not be relied upon. If the $\langle sequence \rangle$ is non-empty, stores the top item from a $\langle sequence \rangle$ in the $\langle token\ list\ variable \rangle$ without removing it from the $\langle sequence \rangle$. The $\langle token\ list\ variable \rangle$ is assigned locally.

`\seq_pop:NNTF`
`\seq_pop:cNTF`
 New: 2012-05-14
 Updated: 2012-05-19

`\seq_pop:NNTF` $\langle sequence \rangle$ $\langle token list variable \rangle$ $\{\langle true code \rangle\}$ $\{\langle false code \rangle\}$

If the $\langle sequence \rangle$ is empty, leaves the $\langle false code \rangle$ in the input stream. The value of the $\langle token list variable \rangle$ is not defined in this case and should not be relied upon. If the $\langle sequence \rangle$ is non-empty, pops the top item from the $\langle sequence \rangle$ in the $\langle token list variable \rangle$, *i.e.* removes the item from the $\langle sequence \rangle$. Both the $\langle sequence \rangle$ and the $\langle token list variable \rangle$ are assigned locally.

`\seq_gpop:NNTF`
`\seq_gpop:cNTF`
 New: 2012-05-14
 Updated: 2012-05-19

`\seq_gpop:NNTF` $\langle sequence \rangle$ $\langle token list variable \rangle$ $\{\langle true code \rangle\}$ $\{\langle false code \rangle\}$

If the $\langle sequence \rangle$ is empty, leaves the $\langle false code \rangle$ in the input stream. The value of the $\langle token list variable \rangle$ is not defined in this case and should not be relied upon. If the $\langle sequence \rangle$ is non-empty, pops the top item from the $\langle sequence \rangle$ in the $\langle token list variable \rangle$, *i.e.* removes the item from the $\langle sequence \rangle$. The $\langle sequence \rangle$ is modified globally, while the $\langle token list variable \rangle$ is assigned locally.

`\seq_push:Nn`
`\seq_push:(NV|Nv|No|Nx|cn|cV|cv|co|cx)`
`\seq_gpush:Nn`
`\seq_gpush:(NV|Nv|No|Nx|cn|cV|cv|co|cx)`

`\seq_push:Nn` $\langle sequence \rangle$ $\{\langle item \rangle\}$

Adds the $\{\langle item \rangle\}$ to the top of the $\langle sequence \rangle$.

10 Sequences as sets

Sequences can also be used as sets, such that all of their items are distinct. Usage of sequences as sets is not currently widespread, hence no specific set function is provided. Instead, it is explained here how common set operations can be performed by combining several functions described in earlier sections. When using sequences to implement sets, one should be careful not to rely on the order of items in the sequence representing the set.

Sets should not contain several occurrences of a given item. To make sure that a $\langle sequence variable \rangle$ only has distinct items, use `\seq_remove_duplicates:N` $\langle sequence variable \rangle$. This function is relatively slow, and to avoid performance issues one should only use it when necessary.

Some operations on a set $\langle seq var \rangle$ are straightforward. For instance, `\seq_count:N` $\langle seq var \rangle$ expands to the number of items, while `\seq_if_in:NnTF` $\langle seq var \rangle$ $\{\langle item \rangle\}$ tests if the $\langle item \rangle$ is in the set.

Adding an $\langle item \rangle$ to a set $\langle seq var \rangle$ can be done by appending it to the $\langle seq var \rangle$ if it is not already in the $\langle seq var \rangle$:

`\seq_if_in:NnF` $\langle seq var \rangle$ $\{\langle item \rangle\}$
`{ \seq_put_right:Nn` $\langle seq var \rangle$ $\{\langle item \rangle\}$ `}`

Removing an $\langle item \rangle$ from a set $\langle seq var \rangle$ can be done using `\seq_remove_all:Nn`,

`\seq_remove_all:Nn` $\langle seq var \rangle$ $\{\langle item \rangle\}$

The intersection of two sets $\langle seq var_1 \rangle$ and $\langle seq var_2 \rangle$ can be stored into $\langle seq var_3 \rangle$ by collecting items of $\langle seq var_1 \rangle$ which are in $\langle seq var_2 \rangle$.

```

\seq_clear:N <seq var3>
\seq_map_inline:Nn <seq var1>
{
\seq_if_in:NnT <seq var2> {#1}
{ \seq_put_right:Nn <seq var3> {#1} }
}

```

The code as written here only works if $\langle seq\ var_3 \rangle$ is different from the other two sequence variables. To cover all cases, items should first be collected in a sequence $\backslash l_ \langle pkg \rangle_internal_seq$, then $\langle seq\ var_3 \rangle$ should be set equal to this internal sequence. The same remark applies to other set functions.

The union of two sets $\langle seq\ var_1 \rangle$ and $\langle seq\ var_2 \rangle$ can be stored into $\langle seq\ var_3 \rangle$ through

```

\seq_concat:NNN <seq var3> <seq var1> <seq var2>
\seq_remove_duplicates:N <seq var3>

```

or by adding items to (a copy of) $\langle seq\ var_1 \rangle$ one by one

```

\seq_set_eq:NN <seq var3> <seq var1>
\seq_map_inline:Nn <seq var2>
{
\seq_if_in:NnF <seq var3> {#1}
{ \seq_put_right:Nn <seq var3> {#1} }
}

```

The second approach is faster than the first when the $\langle seq\ var_2 \rangle$ is short compared to $\langle seq\ var_1 \rangle$.

The difference of two sets $\langle seq\ var_1 \rangle$ and $\langle seq\ var_2 \rangle$ can be stored into $\langle seq\ var_3 \rangle$ by removing items of the $\langle seq\ var_2 \rangle$ from (a copy of) the $\langle seq\ var_1 \rangle$ one by one.

```

\seq_set_eq:NN <seq var3> <seq var1>
\seq_map_inline:Nn <seq var2>
{ \seq_remove_all:Nn <seq var3> {#1} }

```

The symmetric difference of two sets $\langle seq\ var_1 \rangle$ and $\langle seq\ var_2 \rangle$ can be stored into $\langle seq\ var_3 \rangle$ by computing the difference between $\langle seq\ var_1 \rangle$ and $\langle seq\ var_2 \rangle$ and storing the result as $\backslash l_ \langle pkg \rangle_internal_seq$, then the difference between $\langle seq\ var_2 \rangle$ and $\langle seq\ var_1 \rangle$, and finally concatenating the two differences to get the symmetric differences.

```

\seq_set_eq:NN \l\_ \langle pkg \rangle\_internal\_seq <seq var1>
\seq_map_inline:Nn <seq var2>
{ \seq_remove_all:Nn \l\_ \langle pkg \rangle\_internal\_seq {#1} }
\seq_set_eq:NN <seq var3> <seq var2>
\seq_map_inline:Nn <seq var1>
{ \seq_remove_all:Nn <seq var3> {#1} }
\seq_concat:NNN <seq var3> <seq var3> \l\_ \langle pkg \rangle\_internal\_seq

```

11 Constant and scratch sequences

$\backslash c_empty_seq$ Constant that is always empty.

New: 2012-07-02

`\l_tmpa_seq`
`\l_tmpb_seq`

New: 2012-04-26

Scratch sequences for local assignment. These are never used by the kernel code, and so are safe for use with any L^AT_EX3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.

`\g_tmpa_seq`
`\g_tmpb_seq`

New: 2012-04-26

Scratch sequences for global assignment. These are never used by the kernel code, and so are safe for use with any L^AT_EX3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.

12 Viewing sequences

`\seq_show:N`
`\seq_show:c`

Updated: 2015-08-01

`\seq_show:N` $\langle sequence \rangle$
Displays the entries in the $\langle sequence \rangle$ in the terminal.

`\seq_log:N`
`\seq_log:c`

New: 2014-08-12
Updated: 2015-08-01

`\seq_log:N` $\langle sequence \rangle$
Writes the entries in the $\langle sequence \rangle$ in the log file.

Part XI

The l3int package

Integers

Calculation and comparison of integer values can be carried out using literal numbers, `int` registers, constants and integers stored in token list variables. The standard operators `+`, `-`, `/` and `*` and parentheses can be used within such expressions to carry arithmetic operations. This module carries out these functions on *integer expressions* (“`intexpr`”).

1 Integer expressions

`\int_eval:n *` `\int_eval:n {(integer expression)}`

Evaluates the *integer expression* and leaves the result in the input stream as an integer denotation: for positive results an explicit sequence of decimal digits not starting with 0, for negative results - followed by such a sequence, and 0 for zero. The *integer expression* should consist, after expansion, of +, -, *, /, (,) and of course integer operands. The result is calculated by applying standard mathematical rules with the following peculiarities:

- / denotes division rounded to the closest integer with ties rounded away from zero;
- there is an error and the overall expression evaluates to zero whenever the absolute value of any intermediate result exceeds $2^{31} - 1$, except in the case of scaling operations $a*b/c$, for which $a*b$ may be arbitrarily large;
- parentheses may not appear after unary + or -, namely placing +(or -(at the start of an expression or after +, -, *, / or (leads to an error.

Each integer operand can be either an integer variable (with no need for `\int_use:N`) or an integer denotation. For example both

```
\int_eval:n { 5 + 4 * 3 - ( 3 + 4 * 5 ) }
```

and

```
\tl_new:N \l_my_tl
\tl_set:Nn \l_my_tl { 5 }
\int_new:N \l_my_int
\int_set:Nn \l_my_int { 4 }
\int_eval:n { \l_my_tl + \l_my_int * 3 - ( 3 + 4 * 5 ) }
```

evaluate to -6 because `\l_my_tl` expands to the integer denotation 5. As the *integer expression* is fully expanded from left to right during evaluation, fully expandable and restricted-expandable functions can both be used, and `\exp_not:n` and its variants have no effect while `\exp_not:N` may incorrectly interrupt the expression.

T_EXhackers note: Exactly two expansions are needed to evaluate `\int_eval:n`. The result is *not* an *internal integer*, and therefore requires suitable termination if used in a T_EX-style integer assignment.

As all T_EX integers, integer operands can also be dimension or skip variables, converted to integers in `sp`, or octal numbers given as ' followed by digits other than 8 and 9, or hexadecimal numbers given as " followed by digits or upper case letters from A to F, or the character code of some character or one-character control sequence, given as 'char'.

<hr/> <code>\int_eval:w</code> ★ <hr/>	<code>\int_eval:w</code> $\langle integer\ expression \rangle$
New: 2018-03-30	Evaluates the $\langle integer\ expression \rangle$ as described for <code>\int_eval:n</code> . The end of the expression is the first token encountered that cannot form part of such an expression. If that token is <code>\scan_stop</code> : it is removed, otherwise not. Spaces do <i>not</i> terminate the expression. However, spaces terminate explicit integers, and this may terminate the expression: for instance, <code>\int_eval:w 1_+1_9</code> expands to 29 since the digit 9 is not part of the expression.
<hr/> <code>\int_sign:n</code> ★ <hr/>	<code>\int_sign:n</code> $\{\langle intexpr \rangle\}$
New: 2018-11-03	Evaluates the $\langle integer\ expression \rangle$ then leaves 1 or 0 or -1 in the input stream according to the sign of the result.
<hr/> <code>\int_abs:n</code> ★ <hr/>	<code>\int_abs:n</code> $\{\langle integer\ expression \rangle\}$
Updated: 2012-09-26	Evaluates the $\langle integer\ expression \rangle$ as described for <code>\int_eval:n</code> and leaves the absolute value of the result in the input stream as an $\langle integer\ denotation \rangle$ after two expansions.
<hr/> <code>\int_div_round:nn</code> ★ <hr/>	<code>\int_div_round:nn</code> $\{\langle intexpr_1 \rangle\} \{\langle intexpr_2 \rangle\}$
Updated: 2012-09-26	Evaluates the two $\langle integer\ expressions \rangle$ as described earlier, then divides the first value by the second, and rounds the result to the closest integer. Ties are rounded away from zero. Note that this is identical to using <code>/</code> directly in an $\langle integer\ expression \rangle$. The result is left in the input stream as an $\langle integer\ denotation \rangle$ after two expansions.
<hr/> <code>\int_div_truncate:nn</code> ★ <hr/>	<code>\int_div_truncate:nn</code> $\{\langle intexpr_1 \rangle\} \{\langle intexpr_2 \rangle\}$
Updated: 2012-02-09	Evaluates the two $\langle integer\ expressions \rangle$ as described earlier, then divides the first value by the second, and rounds the result towards zero. Note that division using <code>/</code> rounds to the closest integer instead. The result is left in the input stream as an $\langle integer\ denotation \rangle$ after two expansions.
<hr/> <code>\int_max:nn</code> ★ <hr/>	<code>\int_max:nn</code> $\{\langle intexpr_1 \rangle\} \{\langle intexpr_2 \rangle\}$
<code>\int_min:nn</code> ★	<code>\int_min:nn</code> $\{\langle intexpr_1 \rangle\} \{\langle intexpr_2 \rangle\}$
Updated: 2012-09-26	Evaluates the $\langle integer\ expressions \rangle$ as described for <code>\int_eval:n</code> and leaves either the larger or smaller value in the input stream as an $\langle integer\ denotation \rangle$ after two expansions.
<hr/> <code>\int_mod:nn</code> ★ <hr/>	<code>\int_mod:nn</code> $\{\langle intexpr_1 \rangle\} \{\langle intexpr_2 \rangle\}$
Updated: 2012-09-26	Evaluates the two $\langle integer\ expressions \rangle$ as described earlier, then calculates the integer remainder of dividing the first expression by the second. This is obtained by subtracting <code>\int_div_truncate:nn</code> $\{\langle intexpr_1 \rangle\} \{\langle intexpr_2 \rangle\}$ times $\langle intexpr_2 \rangle$ from $\langle intexpr_1 \rangle$. Thus, the result has the same sign as $\langle intexpr_1 \rangle$ and its absolute value is strictly less than that of $\langle intexpr_2 \rangle$. The result is left in the input stream as an $\langle integer\ denotation \rangle$ after two expansions.

2 Creating and initialising integers

<hr/> <code>\int_new:N</code> <hr/>	<code>\int_new:N</code> $\langle integer \rangle$
<code>\int_new:c</code>	Creates a new $\langle integer \rangle$ or raises an error if the name is already taken. The declaration is global. The $\langle integer \rangle$ is initially equal to 0.

<code>\int_const:Nn</code>	<code>\int_const:Nn <integer> {<integer expression>}</code>
<code>\int_const:cn</code>	
Updated: 2011-10-22	Creates a new constant $\langle integer \rangle$ or raises an error if the name is already taken. The value of the $\langle integer \rangle$ is set globally to the $\langle integer expression \rangle$.

<code>\int_zero:N</code>	<code>\int_zero:N <integer></code>
<code>\int_zero:c</code>	
<code>\int_gzero:N</code>	Sets $\langle integer \rangle$ to 0.
<code>\int_gzero:c</code>	

<code>\int_zero_new:N</code>	<code>\int_zero_new:N <integer></code>
<code>\int_zero_new:c</code>	
<code>\int_gzero_new:N</code>	Ensures that the $\langle integer \rangle$ exists globally by applying <code>\int_new:N</code> if necessary, then applies <code>\int_(g)zero:N</code> to leave the $\langle integer \rangle$ set to zero.
<code>\int_gzero_new:c</code>	

New: 2011-12-13

<code>\int_set_eq:NN</code>	<code>\int_set_eq:NN <integer₁₂</code>
<code>\int_set_eq:(cN Nc cc)</code>	
<code>\int_gset_eq:NN</code>	Sets the content of $\langle integer_1 \rangle$ equal to that of $\langle integer_2 \rangle$.
<code>\int_gset_eq:(cN Nc cc)</code>	

<code>\int_if_exist_p:N *</code>	<code>\int_if_exist_p:N <int></code>
<code>\int_if_exist_p:c *</code>	<code>\int_if_exist:NTF <int> {<true code>} {<false code>}</code>
<code>\int_if_exist:NTF *</code>	
<code>\int_if_exist:cTF *</code>	Tests whether the $\langle int \rangle$ is currently defined. This does not check that the $\langle int \rangle$ really is an integer variable.

New: 2012-03-03

3 Setting and incrementing integers

<code>\int_add:Nn</code>	<code>\int_add:Nn <integer> {<integer expression>}</code>
<code>\int_add:cn</code>	
<code>\int_gadd:Nn</code>	Adds the result of the $\langle integer expression \rangle$ to the current content of the $\langle integer \rangle$.
<code>\int_gadd:cn</code>	

Updated: 2011-10-22

<code>\int_decr:N</code>	<code>\int_decr:N <integer></code>
<code>\int_decr:c</code>	
<code>\int_gdecr:N</code>	Decreases the value stored in $\langle integer \rangle$ by 1.
<code>\int_gdecr:c</code>	

<code>\int_incr:N</code>	<code>\int_incr:N <integer></code>
<code>\int_incr:c</code>	
<code>\int_gincr:N</code>	Increases the value stored in $\langle integer \rangle$ by 1.
<code>\int_gincr:c</code>	

<code>\int_set:Nn</code>	<code>\int_set:Nn <integer> {<integer expression>}</code>
<code>\int_set:cn</code>	
<code>\int_gset:Nn</code>	Sets <i><integer></i> to the value of <i><integer expression></i> , which must evaluate to an integer (as described for <code>\int_eval:n</code>).
<code>\int_gset:cn</code>	

Updated: 2011-10-22

<code>\int_sub:Nn</code>	<code>\int_sub:Nn <integer> {<integer expression>}</code>
<code>\int_sub:cn</code>	
<code>\int_gsub:Nn</code>	Subtracts the result of the <i><integer expression></i> from the current content of the <i><integer></i> .
<code>\int_gsub:cn</code>	

Updated: 2011-10-22

4 Using integers

<code>\int_use:N</code> *	<code>\int_use:N <integer></code>
<code>\int_use:c</code> *	
	Recovers the content of an <i><integer></i> and places it directly in the input stream. An error is raised if the variable does not exist or if it is invalid. Can be omitted in places where an <i><integer></i> is required (such as in the first and third arguments of <code>\int_compare:nNnTF</code>).

Updated: 2011-10-22

TeXhackers note: `\int_use:N` is the TeX primitive `\the`: this is one of several L^AT_EX3 names for this primitive.

5 Integer expression conditionals

<code>\int_compare_p:nNn</code> *	<code>\int_compare_p:nNn {<intexpr₁>} <relation> {<intexpr₂>}</code>
<code>\int_compare:nNnTF</code> *	<code>\int_compare:nNnTF {<intexpr₁>} <relation> {<intexpr₂>} {<true code>} {<false code>}</code>

This function first evaluates each of the *<integer expressions>* as described for `\int_eval:n`. The two results are then compared using the *<relation>*:

Equal	=
Greater than	>
Less than	<

This function is less flexible than `\int_compare:nTF` but around 5 times faster.

```

\int_compare_p:n * \int_compare_p:n
\int_compare:nTF * {
    <intexpr1> <relation1>
    ...
    <intexprN> <relationN>
    <intexprN+1>
}
\int_compare:nTF
{
    <intexpr1> <relation1>
    ...
    <intexprN> <relationN>
    <intexprN+1>
}
{<true code>} {<false code>}

```

Updated: 2013-01-13

This function evaluates the *<integer expressions>* as described for `\int_eval:n` and compares consecutive result using the corresponding *<relation>*, namely it compares *<intexpr₁>* and *<intexpr₂>* using the *<relation₁>*, then *<intexpr₂>* and *<intexpr₃>* using the *<relation₂>*, until finally comparing *<intexpr_N>* and *<intexpr_{N+1}>* using the *<relation_N>*. The test yields **true** if all comparisons are **true**. Each *<integer expression>* is evaluated only once, and the evaluation is lazy, in the sense that if one comparison is **false**, then no other *<integer expression>* is evaluated and no other comparison is performed. The *<relations>* can be any of the following:

Equal	= or ==
Greater than or equal to	>=
Greater than	>
Less than or equal to	<=
Less than	<
Not equal	!=

This function is more flexible than `\int_compare:nNnTF` but around 5 times slower.

<code>\int_case:nn</code> *	<code>\int_case:nnTF {⟨test integer expression⟩}</code>
<code>\int_case:nnTF</code> *	{
	{⟨intexpr case ₁ ⟩} {⟨code case ₁ ⟩}
	{⟨intexpr case ₂ ⟩} {⟨code case ₂ ⟩}
	...
	{⟨intexpr case _n ⟩} {⟨code case _n ⟩}
	}
	{⟨true code⟩}
	{⟨false code⟩}

New: 2013-07-24

This function evaluates the *⟨test integer expression⟩* and compares this in turn to each of the *⟨integer expression cases⟩*. If the two are equal then the associated *⟨code⟩* is left in the input stream and other cases are discarded. If any of the cases are matched, the *⟨true code⟩* is also inserted into the input stream (after the code for the appropriate case), while if none match then the *⟨false code⟩* is inserted. The function `\int_case:nn`, which does nothing if there is no match, is also available. For example

```
\int_case:nnF
{ 2 * 5 }
{
  { 5 }      { Small }
  { 4 + 6 }   { Medium }
  { -2 * 10 } { Negative }
}
{ No idea! }
```

leaves “Medium” in the input stream.

<code>\int_if_even_p:n</code> *	<code>\int_if_odd_p:n {⟨integer expression⟩}</code>
<code>\int_if_even:nTF</code> *	<code>\int_if_odd:nTF {⟨integer expression⟩}</code>
<code>\int_if_odd_p:n</code> *	{⟨true code⟩} {⟨false code⟩}
<code>\int_if_odd:nTF</code> *	

This function first evaluates the *⟨integer expression⟩* as described for `\int_eval:n`. It then evaluates if this is odd or even, as appropriate.

6 Integer expression loops

<code>\int_do_until:nNnn</code> ☆	<code>\int_do_until:nNnn {⟨intexpr₁⟩} ⟨relation⟩ {⟨intexpr₂⟩} {⟨code⟩}</code>
-----------------------------------	---

Places the *⟨code⟩* in the input stream for T_EX to process, and then evaluates the relationship between the two *⟨integer expressions⟩* as described for `\int_compare:nNnTF`. If the test is **false** then the *⟨code⟩* is inserted into the input stream again and a loop occurs until the *⟨relation⟩* is **true**.

<code>\int_do_while:nNnn</code> ☆	<code>\int_do_while:nNnn {⟨intexpr₁⟩} ⟨relation⟩ {⟨intexpr₂⟩} {⟨code⟩}</code>
-----------------------------------	---

Places the *⟨code⟩* in the input stream for T_EX to process, and then evaluates the relationship between the two *⟨integer expressions⟩* as described for `\int_compare:nNnTF`. If the test is **true** then the *⟨code⟩* is inserted into the input stream again and a loop occurs until the *⟨relation⟩* is **false**.

<hr/> <code>\int_until_do:nNnn</code> ☆ <hr/>	<code>\int_until_do:nNnn {<intexpr1>} <relation> {<intexpr2>} {<code>}</code>
	Evaluates the relationship between the two <i><integer expressions></i> as described for <code>\int_compare:nNnTF</code> , and then places the <i><code></i> in the input stream if the <i><relation></i> is false . After the <i><code></i> has been processed by T _E X the test is repeated, and a loop occurs until the test is true .
<hr/> <code>\int_while_do:nNnn</code> ☆ <hr/>	<code>\int_while_do:nNnn {<intexpr1>} <relation> {<intexpr2>} {<code>}</code>
	Evaluates the relationship between the two <i><integer expressions></i> as described for <code>\int_compare:nNnTF</code> , and then places the <i><code></i> in the input stream if the <i><relation></i> is true . After the <i><code></i> has been processed by T _E X the test is repeated, and a loop occurs until the test is false .
<hr/> <code>\int_do_until:nn</code> ☆ <hr/>	<code>\int_do_until:nn {<integer relation>} {<code>}</code>
Updated: 2013-01-13	Places the <i><code></i> in the input stream for T _E X to process, and then evaluates the <i><integer relation></i> as described for <code>\int_compare:nTF</code> . If the test is false then the <i><code></i> is inserted into the input stream again and a loop occurs until the <i><relation></i> is true .
<hr/> <code>\int_do_while:nn</code> ☆ <hr/>	<code>\int_do_while:nn {<integer relation>} {<code>}</code>
Updated: 2013-01-13	Places the <i><code></i> in the input stream for T _E X to process, and then evaluates the <i><integer relation></i> as described for <code>\int_compare:nTF</code> . If the test is true then the <i><code></i> is inserted into the input stream again and a loop occurs until the <i><relation></i> is false .
<hr/> <code>\int_until_do:nn</code> ☆ <hr/>	<code>\int_until_do:nn {<integer relation>} {<code>}</code>
Updated: 2013-01-13	Evaluates the <i><integer relation></i> as described for <code>\int_compare:nTF</code> , and then places the <i><code></i> in the input stream if the <i><relation></i> is false . After the <i><code></i> has been processed by T _E X the test is repeated, and a loop occurs until the test is true .
<hr/> <code>\int_while_do:nn</code> ☆ <hr/>	<code>\int_while_do:nn {<integer relation>} {<code>}</code>
Updated: 2013-01-13	Evaluates the <i><integer relation></i> as described for <code>\int_compare:nTF</code> , and then places the <i><code></i> in the input stream if the <i><relation></i> is true . After the <i><code></i> has been processed by T _E X the test is repeated, and a loop occurs until the test is false .

7 Integer step functions

<code>\int_step_function:nN</code>	☆	<code>\int_step_function:nN {⟨final value⟩} ⟨function⟩</code>
<code>\int_step_function:nnN</code>	☆	<code>\int_step_function:nnN {⟨initial value⟩} {⟨final value⟩} ⟨function⟩</code>
<code>\int_step_function:nnnN</code>	☆	<code>\int_step_function:nnnN {⟨initial value⟩} {⟨step⟩} {⟨final value⟩} ⟨function⟩</code>

New: 2012-06-04
Updated: 2018-04-22

This function first evaluates the $\langle initial\ value \rangle$, $\langle step \rangle$ and $\langle final\ value \rangle$, all of which should be integer expressions. The $\langle function \rangle$ is then placed in front of each $\langle value \rangle$ from the $\langle initial\ value \rangle$ to the $\langle final\ value \rangle$ in turn (using $\langle step \rangle$ between each $\langle value \rangle$). The $\langle step \rangle$ must be non-zero. If the $\langle step \rangle$ is positive, the loop stops when the $\langle value \rangle$ becomes larger than the $\langle final\ value \rangle$. If the $\langle step \rangle$ is negative, the loop stops when the $\langle value \rangle$ becomes smaller than the $\langle final\ value \rangle$. The $\langle function \rangle$ should absorb one numerical argument. For example

```
\cs_set:Npn \my_func:n #1 { [I~saw~#1] \quad }
\int_step_function:nnnN { 1 } { 1 } { 5 } \my_func:n
```

would print

```
[I saw 1]   [I saw 2]   [I saw 3]   [I saw 4]   [I saw 5]
```

The functions `\int_step_function:nN` and `\int_step_function:nnN` both use a fixed $\langle step \rangle$ of 1, and in the case of `\int_step_function:nN` the $\langle initial\ value \rangle$ is also fixed as 1. These functions are provided as simple short-cuts for code clarity.

<code>\int_step_inline:nn</code>	<code>\int_step_inline:nn {⟨final value⟩} {⟨code⟩}</code>
<code>\int_step_inline:nnn</code>	<code>\int_step_inline:nnn {⟨initial value⟩} {⟨final value⟩} {⟨code⟩}</code>
<code>\int_step_inline:nnnn</code>	<code>\int_step_inline:nnnn {⟨initial value⟩} {⟨step⟩} {⟨final value⟩} {⟨code⟩}</code>

New: 2012-06-04
Updated: 2018-04-22

This function first evaluates the $\langle initial\ value \rangle$, $\langle step \rangle$ and $\langle final\ value \rangle$, all of which should be integer expressions. Then for each $\langle value \rangle$ from the $\langle initial\ value \rangle$ to the $\langle final\ value \rangle$ in turn (using $\langle step \rangle$ between each $\langle value \rangle$), the $\langle code \rangle$ is inserted into the input stream with `#1` replaced by the current $\langle value \rangle$. Thus the $\langle code \rangle$ should define a function of one argument (`#1`).

The functions `\int_step_inline:nn` and `\int_step_inline:nnn` both use a fixed $\langle step \rangle$ of 1, and in the case of `\int_step_inline:nn` the $\langle initial\ value \rangle$ is also fixed as 1. These functions are provided as simple short-cuts for code clarity.

<code>\int_step_variable:nNn</code>	<code>\int_step_variable:nNn {⟨final value⟩} ⟨tl var⟩ {⟨code⟩}</code>
<code>\int_step_variable:nnNn</code>	<code>\int_step_variable:nnNn {⟨initial value⟩} {⟨final value⟩} ⟨tl var⟩ {⟨code⟩}</code>
<code>\int_step_variable:nnnNn</code>	<code>\int_step_variable:nnnNn {⟨initial value⟩} {⟨step⟩} {⟨final value⟩} ⟨tl var⟩ {⟨code⟩}</code>

New: 2012-06-04
Updated: 2018-04-22

This function first evaluates the $\langle initial\ value \rangle$, $\langle step \rangle$ and $\langle final\ value \rangle$, all of which should be integer expressions. Then for each $\langle value \rangle$ from the $\langle initial\ value \rangle$ to the $\langle final\ value \rangle$ in turn (using $\langle step \rangle$ between each $\langle value \rangle$), the $\langle code \rangle$ is inserted into the input stream, with the $\langle tl\ var \rangle$ defined as the current $\langle value \rangle$. Thus the $\langle code \rangle$ should make use of the $\langle tl\ var \rangle$.

The functions `\int_step_variable:nNn` and `\int_step_variable:nnNn` both use a fixed $\langle step \rangle$ of 1, and in the case of `\int_step_variable:nNn` the $\langle initial\ value \rangle$ is also fixed as 1. These functions are provided as simple short-cuts for code clarity.

8 Formatting integers

Integers can be placed into the output stream with formatting. These conversions apply to any integer expressions.

<code>\int_to_arabic:n *</code>	<code>\int_to_arabic:n {⟨integer expression⟩}</code>
---------------------------------	--

Updated: 2011-10-22

Places the value of the *⟨integer expression⟩* in the input stream as digits, with category code 12 (other).

<code>\int_to_alph:n *</code> <code>\int_to_Alph:n *</code>	<code>\int_to_alph:n {⟨integer expression⟩}</code>
--	--

Updated: 2011-09-17

Evaluates the *⟨integer expression⟩* and converts the result into a series of letters, which are then left in the input stream. The conversion rule uses the 26 letters of the English alphabet, in order, adding letters when necessary to increase the total possible range of representable numbers. Thus

```
\int_to_alph:n { 1 }
```

places a in the input stream,

```
\int_to_alph:n { 26 }
```

is represented as z and

```
\int_to_alph:n { 27 }
```

is converted to aa. For conversions using other alphabets, use `\int_to_symbols:nnn` to define an alphabet-specific function. The basic `\int_to_alph:n` and `\int_to_Alph:n` functions should not be modified. The resulting tokens are digits with category code 12 (other) and letters with category code 11 (letter).

<code>\int_to_symbols:nnn *</code>	<code>\int_to_symbols:nnn</code> <code>{⟨integer expression⟩} {⟨total symbols⟩}</code> <code>{⟨value to symbol mapping⟩}</code>
------------------------------------	---

Updated: 2011-09-17

This is the low-level function for conversion of an *⟨integer expression⟩* into a symbolic form (often letters). The *⟨total symbols⟩* available should be given as an integer expression. Values are actually converted to symbols according to the *⟨value to symbol mapping⟩*. This should be given as *⟨total symbols⟩* pairs of entries, a number and the appropriate symbol. Thus the `\int_to_alph:n` function is defined as

```
\cs_new:Npn \int_to_alph:n #1
{
  \int_to_symbols:nnn {#1} { 26 }
  {
    { 1 } { a }
    { 2 } { b }
    ...
    { 26 } { z }
  }
}
```

<hr/>	
<code>\int_to_bin:n *</code>	<code>\int_to_bin:n {⟨integer expression⟩}</code>
<hr/>	
New: 2014-02-11	Calculates the value of the $\langle integer\ expression \rangle$ and places the binary representation of the result in the input stream.
<hr/>	
<code>\int_to_hex:n *</code>	<code>\int_to_hex:n {⟨integer expression⟩}</code>
<code>\int_to_Hex:n *</code>	
<hr/>	
New: 2014-02-11	Calculates the value of the $\langle integer\ expression \rangle$ and places the hexadecimal (base 16) representation of the result in the input stream. Letters are used for digits beyond 9: lower case letters for <code>\int_to_hex:n</code> and upper case ones for <code>\int_to_Hex:n</code> . The resulting tokens are digits with category code 12 (other) and letters with category code 11 (letter).
<hr/>	
<code>\int_to_oct:n *</code>	<code>\int_to_oct:n {⟨integer expression⟩}</code>
<hr/>	
New: 2014-02-11	Calculates the value of the $\langle integer\ expression \rangle$ and places the octal (base 8) representation of the result in the input stream. The resulting tokens are digits with category code 12 (other) and letters with category code 11 (letter).
<hr/>	
<code>\int_to_base:nn *</code>	<code>\int_to_base:nn {⟨integer expression⟩} {⟨base⟩}</code>
<code>\int_to_Base:nn *</code>	
<hr/>	
Updated: 2014-02-11	Calculates the value of the $\langle integer\ expression \rangle$ and converts it into the appropriate representation in the $\langle base \rangle$; the later may be given as an integer expression. For bases greater than 10 the higher “digits” are represented by letters from the English alphabet: lower case letters for <code>\int_to_base:n</code> and upper case ones for <code>\int_to_Base:n</code> . The maximum $\langle base \rangle$ value is 36. The resulting tokens are digits with category code 12 (other) and letters with category code 11 (letter).
<hr/>	
TeXhackers note: This is a generic version of <code>\int_to_bin:n</code> , etc.	
<hr/>	
<code>\int_to_roman:n ☆</code>	<code>\int_to_roman:n {⟨integer expression⟩}</code>
<code>\int_to_Roman:n ☆</code>	
<hr/>	
Updated: 2011-10-22	Places the value of the $\langle integer\ expression \rangle$ in the input stream as Roman numerals, either lower case (<code>\int_to_roman:n</code>) or upper case (<code>\int_to_Roman:n</code>). If the value is negative or zero, the output is empty. The Roman numerals are letters with category code 11 (letter). The letters used are <code>mdclxvi</code> , repeated as needed: the notation with bars (such as \bar{v} for 5000) is <i>not</i> used. For instance <code>\int_to_roman:n { 8249 }</code> expands to <code>mmmmmmmmccxlix</code> .

9 Converting from other formats to integers

<hr/>	
<code>\int_from_alph:n *</code>	<code>\int_from_alph:n {⟨letters⟩}</code>
<hr/>	
Updated: 2014-08-25	Converts the $\langle letters \rangle$ into the integer (base 10) representation and leaves this in the input stream. The $\langle letters \rangle$ are first converted to a string, with no expansion. Lower and upper case letters from the English alphabet may be used, with “a” equal to 1 through to “z” equal to 26. The function also accepts a leading sign, made of + and -. This is the inverse function of <code>\int_to_alph:n</code> and <code>\int_to_Alph:n</code> .

<hr/> <code>\int_from_bin:n</code> ★ <hr/>	<code>\int_from_bin:n {⟨binary number⟩}</code>
New: 2014-02-11 Updated: 2014-08-25 <hr/>	Converts the <i>⟨binary number⟩</i> into the integer (base 10) representation and leaves this in the input stream. The <i>⟨binary number⟩</i> is first converted to a string, with no expansion. The function accepts a leading sign, made of + and -, followed by binary digits. This is the inverse function of <code>\int_to_bin:n</code> .
<hr/> <code>\int_from_hex:n</code> ★ <hr/>	<code>\int_from_hex:n {⟨hexadecimal number⟩}</code>
New: 2014-02-11 Updated: 2014-08-25 <hr/>	Converts the <i>⟨hexadecimal number⟩</i> into the integer (base 10) representation and leaves this in the input stream. Digits greater than 9 may be represented in the <i>⟨hexadecimal number⟩</i> by upper or lower case letters. The <i>⟨hexadecimal number⟩</i> is first converted to a string, with no expansion. The function also accepts a leading sign, made of + and -. This is the inverse function of <code>\int_to_hex:n</code> and <code>\int_to_Hex:n</code> .
<hr/> <code>\int_from_oct:n</code> ★ <hr/>	<code>\int_from_oct:n {⟨octal number⟩}</code>
New: 2014-02-11 Updated: 2014-08-25 <hr/>	Converts the <i>⟨octal number⟩</i> into the integer (base 10) representation and leaves this in the input stream. The <i>⟨octal number⟩</i> is first converted to a string, with no expansion. The function accepts a leading sign, made of + and -, followed by octal digits. This is the inverse function of <code>\int_to_oct:n</code> .
<hr/> <code>\int_from_roman:n</code> ★ <hr/>	<code>\int_from_roman:n {⟨roman numeral⟩}</code>
Updated: 2014-08-25 <hr/>	Converts the <i>⟨roman numeral⟩</i> into the integer (base 10) representation and leaves this in the input stream. The <i>⟨roman numeral⟩</i> is first converted to a string, with no expansion. The <i>⟨roman numeral⟩</i> may be in upper or lower case; if the numeral contains characters besides <code>mdclxvi</code> or <code>MDCLXVI</code> then the resulting value is -1. This is the inverse function of <code>\int_to_roman:n</code> and <code>\int_to_Roman:n</code> .
<hr/> <code>\int_from_base:nn</code> ★ <hr/>	<code>\int_from_base:nn {⟨number⟩} {⟨base⟩}</code>
Updated: 2014-08-25 <hr/>	Converts the <i>⟨number⟩</i> expressed in <i>⟨base⟩</i> into the appropriate value in base 10. The <i>⟨number⟩</i> is first converted to a string, with no expansion. The <i>⟨number⟩</i> should consist of digits and letters (either lower or upper case), plus optionally a leading sign. The maximum <i>⟨base⟩</i> value is 36. This is the inverse function of <code>\int_to_base:nn</code> and <code>\int_to_Base:nn</code> .

10 Random integers

<hr/> <code>\int_rand:nn</code> ★ <hr/>	<code>\int_rand:nn {⟨integer expr₁⟩} {⟨integer expr₂⟩}</code>
New: 2016-12-06 Updated: 2018-04-27 <hr/>	Evaluates the two <i>⟨integer expressions⟩</i> and produces a pseudo-random number between the two (with bounds included). This is not available in older versions of X _Y TeX.
<hr/> <code>\int_rand:n</code> ★ <hr/>	<code>\int_rand:n {⟨integer expr⟩}</code>
New: 2018-05-05 <hr/>	Evaluates the <i>⟨integer expression⟩</i> then produces a pseudo-random number between 1 and the <i>⟨integer⟩</i> (included). This is not available in older versions of X _Y TeX.

11 Viewing integers

<hr/> <code>\int_show:N</code> <code>\int_show:c</code> <hr/>	<code>\int_show:N <integer></code> Displays the value of the <i><integer></i> on the terminal.
<hr/> <code>\int_show:n</code> <hr/> <div>New: 2011-11-22 Updated: 2015-08-07</div>	<code>\int_show:n {(integer expression)}</code> Displays the result of evaluating the <i><integer expression></i> on the terminal.
<hr/> <code>\int_log:N</code> <code>\int_log:c</code> <hr/> <div>New: 2014-08-22 Updated: 2015-08-03</div>	<code>\int_log:N <integer></code> Writes the value of the <i><integer></i> in the log file.
<hr/> <code>\int_log:n</code> <hr/> <div>New: 2014-08-22 Updated: 2015-08-07</div>	<code>\int_log:n {(integer expression)}</code> Writes the result of evaluating the <i><integer expression></i> in the log file.

12 Constant integers

<hr/> <code>\c_zero_int</code> <code>\c_one_int</code> <hr/> <div>New: 2018-05-07</div>	Integer values used with primitive tests and assignments: their self-terminating nature makes these more convenient and faster than literal numbers.
<hr/> <code>\c_max_int</code> <hr/>	The maximum value that can be stored as an integer.
<hr/> <code>\c_max_register_int</code> <hr/>	Maximum number of registers.
<hr/> <code>\c_max_char_int</code> <hr/>	Maximum character code completely supported by the engine.

13 Scratch integers

<hr/> <code>\l_tmpa_int</code> <code>\l_tmpb_int</code> <hr/>	Scratch integer for local assignment. These are never used by the kernel code, and so are safe for use with any L ^A T _E X3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.
<hr/> <code>\g_tmpa_int</code> <code>\g_tmpb_int</code> <hr/>	Scratch integer for global assignment. These are never used by the kernel code, and so are safe for use with any L ^A T _E X3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.

14 Direct number expansion

<code>\int_value:w</code>	★
New: 2018-03-27	

`\int_value:w` $\langle integer \rangle$
`\int_value:w` $\langle integer\ denotation \rangle$ $\langle optional\ space \rangle$

Expands the following tokens until an $\langle integer \rangle$ is formed, and leaves a normalized form (no leading sign except for negative numbers, no leading digit 0 except for zero) in the input stream as category code 12 (other) characters. The $\langle integer \rangle$ can consist of any number of signs (with intervening spaces) followed by

- an integer variable (in fact, any T_EX register except `\toks`) or
- explicit digits (or by ‘ $\langle octal\ digits \rangle$ ’ or “ $\langle hexadecimal\ digits \rangle$ ” or ‘ $\langle character \rangle$ ’).

In this last case expansion stops once a non-digit is found; if that is a space it is removed as in `f`-expansion, and so `\exp_stop_f:` may be employed as an end marker. Note that protected functions *are* expanded by this process.

This function requires exactly one expansion to produce a value, and so is suitable for use in cases where a number is required “directly”. In general, `\int_eval:n` is the preferred approach to generating numbers.

T_EXhackers note: This is the T_EX primitive `\number`.

15 Primitive conditionals

<code>\if_int_compare:w</code>	★
--------------------------------	---

`\if_int_compare:w` $\langle integer_1 \rangle$ $\langle relation \rangle$ $\langle integer_2 \rangle$
 $\langle true\ code \rangle$
`\else:`
 $\langle false\ code \rangle$
`\fi:`

Compare two integers using $\langle relation \rangle$, which must be one of =, < or > with category code 12. The `\else:` branch is optional.

T_EXhackers note: These are both names for the T_EX primitive `\ifnum`.

<code>\if_case:w</code>	★
<code>\or:</code>	★

`\if_case:w` $\langle integer \rangle$ $\langle case_0 \rangle$
`\or:` $\langle case_1 \rangle$
`\or:` ...
`\else:` $\langle default \rangle$
`\fi:`

Selects a case to execute based on the value of the $\langle integer \rangle$. The first case ($\langle case_0 \rangle$) is executed if $\langle integer \rangle$ is 0, the second ($\langle case_1 \rangle$) if the $\langle integer \rangle$ is 1, *etc.* The $\langle integer \rangle$ may be a literal, a constant or an integer expression (*e.g.* using `\int_eval:n`).

T_EXhackers note: These are the T_EX primitives `\ifcase` and `\or`.

<u><code>\if_int_odd:w</code></u> *	<code>\if_int_odd:w</code> $\langle tokens \rangle$ $\langle optional\ space \rangle$ $\langle true\ code \rangle$ <code>\else:</code> $\langle true\ code \rangle$ <code>\fi:</code>
-------------------------------------	---

Expands $\langle tokens \rangle$ until a non-numeric token or a space is found, and tests whether the resulting $\langle integer \rangle$ is odd. If so, $\langle true\ code \rangle$ is executed. The `\else:` branch is optional.

TeXhackers note: This is the TeX primitive `\ifodd`.

Part XII

The l3flag package: Expandable flags

Flags are the only data-type that can be modified in expansion-only contexts. This module is meant mostly for kernel use: in almost all cases, booleans or integers should be preferred to flags because they are very significantly faster.

A flag can hold any non-negative value, which we call its *height*. In expansion-only contexts, a flag can only be “raised”: this increases the *height* by 1. The *height* can also be queried expandably. However, decreasing it, or setting it to zero requires non-expandable assignments.

Flag variables are always local. They are referenced by a *flag name* such as `str_missing`. The *flag name* is used as part of `\use:c` constructions hence is expanded at point of use. It must expand to character tokens only, with no spaces.

A typical use case of flags would be to keep track of whether an exceptional condition has occurred during expandable processing, and produce a meaningful (non-expandable) message after the end of the expandable processing. This is exemplified by `l3str-convert`, which for performance reasons performs conversions of individual characters expandably and for readability reasons produces a single error message describing incorrect inputs that were encountered.

Flags should not be used without carefully considering the fact that raising a flag takes a time and memory proportional to its height. Flags should not be used unless unavoidable.

1 Setting up flags

<code>\flag_new:n</code>	<code>\flag_new:n {<flag name>}</code>
--------------------------	--

Creates a new flag with a name given by *flag name*, or raises an error if the name is already taken. The *flag name* may not contain spaces. The declaration is global, but flags are always local variables. The *flag* initially has zero height.

<code>\flag_clear:n</code>	<code>\flag_clear:n {<flag name>}</code>
----------------------------	--

The *flag*’s height is set to zero. The assignment is local.

<code>\flag_clear_new:n</code>	<code>\flag_clear_new:n {<flag name>}</code>
--------------------------------	--

Ensures that the *flag* exists globally by applying `\flag_new:n` if necessary, then applies `\flag_clear:n`, setting the height to zero locally.

<code>\flag_show:n</code>	<code>\flag_show:n {<flag name>}</code>
---------------------------	---

Displays the *flag*’s height in the terminal.

<code>\flag_log:n</code>	<code>\flag_log:n {<flag name>}</code>
--------------------------	--

Writes the *flag*’s height to the log file.

2 Expandable flag commands

<hr/> <code>\flag_if_exist:n *</code> <hr/>	<code>\flag_if_exist:n {⟨flag name⟩}</code>
<code>\flag_if_exist:n\underline{TF} *</code>	This function returns <code>true</code> if the $\langle flag\ name \rangle$ references a flag that has been defined previously, and <code>false</code> otherwise.
<hr/> <code>\flag_if_raised_p:n *</code> <hr/>	<code>\flag_if_raised:n {⟨flag name⟩}</code>
<code>\flag_if_raised:n\underline{TF} *</code>	This function returns <code>true</code> if the $\langle flag \rangle$ has non-zero height, and <code>false</code> if the $\langle flag \rangle$ has zero height.
<hr/> <code>\flag_height:n *</code> <hr/>	<code>\flag_height:n {⟨flag name⟩}</code>
	Expands to the height of the $\langle flag \rangle$ as an integer denotation.
<hr/> <code>\flag_raise:n *</code> <hr/>	<code>\flag_raise:n {⟨flag name⟩}</code>
	The $\langle flag \rangle$'s height is increased by 1 locally.

Part XIII

The l3prg package

Control structures

Conditional processing in L^AT_EX3 is defined as something that performs a series of tests, possibly involving assignments and calling other functions that do not read further ahead in the input stream. After processing the input, a *state* is returned. The states returned are *⟨true⟩* and *⟨false⟩*.

L^AT_EX3 has two forms of conditional flow processing based on these states. The first form is predicate functions that turn the returned state into a boolean *⟨true⟩* or *⟨false⟩*. For example, the function `\cs_if_free_p:N` checks whether the control sequence given as its argument is free and then returns the boolean *⟨true⟩* or *⟨false⟩* values to be used in testing with `\if_predicate:w` or in functions to be described below. The second form is the kind of functions choosing a particular argument from the input stream based on the result of the testing as in `\cs_if_free:NTF` which also takes one argument (the N) and then executes either **true** or **false** depending on the result.

T_EXhackers note: The arguments are executed after exiting the underlying `\if... \fi` structure.

1 Defining a set of conditional functions

```
\prg_new_conditional:Npnn
\prg_set_conditional:Npnn
\prg_new_conditional:Nnn
\prg_set_conditional:Nnn
```

Updated: 2012-02-06

```
\prg_new_conditional:Npnn \<name>:<arg spec> <parameters> {\<conditions>} {\<code>}
\prg_new_conditional:Nnn \<name>:<arg spec> {\<conditions>} {\<code>}
```

These functions create a family of conditionals using the same *{⟨code⟩}* to perform the test created. Those conditionals are expandable if *⟨code⟩* is. The **new** versions check for existing definitions and perform assignments globally (cf. `\cs_new:Npn`) whereas the **set** versions do no check and perform assignments locally (cf. `\cs_set:Npn`). The conditionals created are dependent on the comma-separated list of *⟨conditions⟩*, which should be one or more of p, T, F and TF.

```
\prg_new_protected_conditional:Npnn \prg_new_protected_conditional:Npnn \<name>:<arg spec> <parameters>
\prg_set_protected_conditional:Npnn {\<conditions>} {\<code>}
\prg_new_protected_conditional:Nnn \prg_new_protected_conditional:Nnn \<name>:<arg spec>
\prg_set_protected_conditional:Nnn {\<conditions>} {\<code>}
```

Updated: 2012-02-06

These functions create a family of protected conditionals using the same *{⟨code⟩}* to perform the test created. The *⟨code⟩* does not need to be expandable. The **new** version check for existing definitions and perform assignments globally (cf. `\cs_new:Npn`) whereas the **set** version do not (cf. `\cs_set:Npn`). The conditionals created are depended on the comma-separated list of *⟨conditions⟩*, which should be one or more of T, F and TF (not p).

The conditionals are defined by `\prg_new_conditional:Npnn` and friends as:

- `\<name>_p:<arg spec>` — a predicate function which will supply either a logical `true` or logical `false`. This function is intended for use in cases where one or more logical tests are combined to lead to a final outcome. This function cannot be defined for `protected` conditionals.
- `\<name>:<arg spec>T` — a function with one more argument than the original `<arg spec>` demands. The `<true branch>` code in this additional argument will be left on the input stream only if the test is `true`.
- `\<name>:<arg spec>F` — a function with one more argument than the original `<arg spec>` demands. The `<false branch>` code in this additional argument will be left on the input stream only if the test is `false`.
- `\<name>:<arg spec>TF` — a function with two more argument than the original `<arg spec>` demands. The `<true branch>` code in the first additional argument will be left on the input stream if the test is `true`, while the `<false branch>` code in the second argument will be left on the input stream if the test is `false`.

The `<code>` of the test may use `<parameters>` as specified by the second argument to `\prg_set_conditional:Npnn`: this should match the `<argument specification>` but this is not enforced. The `Nnn` versions infer the number of arguments from the argument specification given (cf. `\cs_new:Nn`, etc.). Within the `<code>`, the functions `\prg_return_true:` and `\prg_return_false:` are used to indicate the logical outcomes of the test.

An example can easily clarify matters here:

```
\prg_set_conditional:Npnn \foo_if_bar:NN #1#2 { p , T , TF }
{
  \if_meaning:w \l_tmpa_tl #1
  \prg_return_true:
\else:
  \if_meaning:w \l_tmpa_tl #2
  \prg_return_true:
\else:
  \prg_return_false:
\fi:
\fi:
}
```

This defines the function `\foo_if_bar_p:NN`, `\foo_if_bar:NNTF` and `\foo_if_bar:NNT` but not `\foo_if_bar:NNF` (because `F` is missing from the `<conditions>` list). The return statements take care of resolving the remaining `\else:` and `\fi:` before returning the state. There must be a return statement for each branch; failing to do so will result in erroneous output if that branch is executed.

<code>\prg_new_eq_conditional:Nnn</code>	<code>\prg_new_eq_conditional:Nnn \<name1>:<arg spec1> \<name2>:<arg spec2></code>
<code>\prg_set_eq_conditional:Nnn</code>	<code>{<conditions>}</code>

These functions copy a family of conditionals. The `new` version checks for existing definitions (cf. `\cs_new_eq:NN`) whereas the `set` version does not (cf. `\cs_set_eq:NN`). The conditionals copied are depended on the comma-separated list of `<conditions>`, which should be one or more of `p`, `T`, `F` and `TF`.

<code>\prg_return_true:</code>	<code>*</code>	<code>\prg_return_true:</code>
<code>\prg_return_false:</code>	<code>*</code>	<code>\prg_return_false:</code>

These “return” functions define the logical state of a conditional statement. They appear within the code for a conditional function generated by `\prg_set_conditional:Npnn`, *etc.*, to indicate when a true or false branch should be taken. While they may appear multiple times each within the code of such conditionals, the execution of the conditional must result in the expansion of one of these two functions *exactly once*.

The return functions trigger what is internally an *f*-expansion process to complete the evaluation of the conditional. Therefore, after `\prg_return_true:` or `\prg_return_false:` there must be no non-expandable material in the input stream for the remainder of the expansion of the conditional code. This includes other instances of either of these functions.

<code>\prg_generate_conditional_variant:Nnn</code>	<code>\prg_generate_conditional_variant:Nnn \<name>:\<arg spec></code>
	<code>{\<variant argument specifiers>} {\<condition specifiers>}</code>

New: 2017-12-12

Defines argument-specifier variants of conditionals. This is equivalent to running `\cs_generate_variant:Nn \<conditional> {\<variant argument specifiers>}` on each *<conditional>* described by the *<condition specifiers>*. These base-form *<conditionals>* are obtained from the *<name>* and *<arg spec>* as described for `\prg_new_conditional:Npnn`, and they should be defined.

2 The boolean data type

This section describes a boolean data type which is closely connected to conditional processing as sometimes you want to execute some code depending on the value of a switch (*e.g.*, draft/final) and other times you perhaps want to use it as a predicate function in an `\if_predicate:w` test. The problem of the primitive `\if_false:` and `\if_true:` tokens is that it is not always safe to pass them around as they may interfere with scanning for termination of primitive conditional processing. Therefore, we employ two canonical booleans: `\c_true_bool` or `\c_false_bool`. Besides preventing problems as described above, it also allows us to implement a simple boolean parser supporting the logical operations And, Or, Not, *etc.* which can then be used on both the boolean type and predicate functions.

All conditional `\bool_` functions except assignments are expandable and expect the input to also be fully expandable (which generally means being constructed from predicate functions and booleans, possibly nested).

T_EXhackers note: The `bool` data type is not implemented using the `\iffalse/\iftrue` primitives, in contrast to `\newif`, *etc.*, in plain T_EX, L^AT_EX 2_ε and so on. Programmers should not base use of `bool` switches on any particular expectation of the implementation.

<code>\bool_new:N</code>	<code>\bool_new:N \<boolean></code>
<code>\bool_new:c</code>	

Creates a new *<boolean>* or raises an error if the name is already taken. The declaration is global. The *<boolean>* is initially **false**.

<hr/> \bool_const:Nn \bool_const:cn <hr/> New: 2017-11-28	\bool_const:Nn <boolean> {<boolexpr>} Creates a new constant <boolean> or raises an error if the name is already taken. The value of the <boolean> is set globally to the result of evaluating the <boolexpr>.
<hr/> \bool_set_false:N \bool_set_false:c \bool_gset_false:N \bool_gset_false:c <hr/>	\bool_set_false:N <boolean> Sets <boolean> logically false.
<hr/> \bool_set_true:N \bool_set_true:c \bool_gset_true:N \bool_gset_true:c <hr/>	\bool_set_true:N <boolean> Sets <boolean> logically true.
<hr/> \bool_set_eq:NN \bool_set_eq:(cN Nc cc) \bool_gset_eq:NN \bool_gset_eq:(cN Nc cc) <hr/>	\bool_set_eq:NN <boolean ₁ > <boolean ₂ > Sets <boolean ₁ > to the current value of <boolean ₂ >.
<hr/> \bool_set:Nn \bool_set:cn \bool_gset:Nn \bool_gset:cn <hr/> Updated: 2017-07-15	\bool_set:Nn <boolean> {<boolexpr>} Evaluates the <boolean expression> as described for \bool_if:nTF, and sets the <boolean> variable to the logical truth of this evaluation.
<hr/> \bool_if_p:N ★ \bool_if_p:c ★ \bool_if:nTF ★ \bool_if:cTF ★ <hr/> Updated: 2017-07-15	\bool_if_p:N <boolean> \bool_if:NTF <boolean> {<true code>} {<false code>} Tests the current truth of <boolean>, and continues expansion based on this result.
<hr/> \bool_show:N \bool_show:c <hr/> New: 2012-02-09 Updated: 2015-08-01	\bool_show:N <boolean> Displays the logical truth of the <boolean> on the terminal.
<hr/> \bool_show:n <hr/> New: 2012-02-09 Updated: 2017-07-15	\bool_show:n {<boolean expression>} Displays the logical truth of the <boolean expression> on the terminal.
<hr/> \bool_log:N \bool_log:c <hr/> New: 2014-08-22 Updated: 2015-08-03	\bool_log:N <boolean> Writes the logical truth of the <boolean> in the log file.

<code>\bool_log:n</code>	<code>\bool_log:n {(boolean expression)}</code>
New: 2014-08-22 Updated: 2017-07-15	Writes the logical truth of the $\langle boolean\ expression \rangle$ in the log file.

<code>\bool_if_exist_p:N *</code>	<code>\bool_if_exist_p:N $\langle boolean \rangle$</code>
<code>\bool_if_exist_p:c *</code>	<code>\bool_if_exist:NTF $\langle boolean \rangle$ {(true code)} {(false code)}</code>
<code>\bool_if_exist:NTF *</code>	Tests whether the $\langle boolean \rangle$ is currently defined. This does not check that the $\langle boolean \rangle$ really is a boolean variable.
<code>\bool_if_exist:cTF *</code>	
New: 2012-03-03	

2.1 Scratch booleans

<code>\l_tmpa_bool</code> <code>\l_tmpb_bool</code>	A scratch boolean for local assignment. It is never used by the kernel code, and so is safe for use with any L ^A T _E X3-defined function. However, it may be overwritten by other non-kernel code and so should only be used for short-term storage.
--	--

<code>\g_tmpa_bool</code> <code>\g_tmpb_bool</code>	A scratch boolean for global assignment. It is never used by the kernel code, and so is safe for use with any L ^A T _E X3-defined function. However, it may be overwritten by other non-kernel code and so should only be used for short-term storage.
--	---

3 Boolean expressions

As we have a boolean datatype and predicate functions returning boolean $\langle true \rangle$ or $\langle false \rangle$ values, it seems only fitting that we also provide a parser for $\langle boolean\ expressions \rangle$.

A boolean expression is an expression which given input in the form of predicate functions and boolean variables, return boolean $\langle true \rangle$ or $\langle false \rangle$. It supports the logical operations And, Or and Not as the well-known infix operators `&&` and `||` and prefix `!` with their usual precedences (namely, `&&` binds more tightly than `||`). In addition to this, parentheses can be used to isolate sub-expressions. For example,

```
\int_compare_p:n { 1 = 1 } &&
(
  \int_compare_p:n { 2 = 3 } ||
  \int_compare_p:n { 4 <= 4 } ||
  \str_if_eq_p:nn { abc } { def }
) &&
! \int_compare_p:n { 2 = 4 }
```

is a valid boolean expression.

Contrarily to some other programming languages, the operators `&&` and `||` evaluate both operands in all cases, even when the first operand is enough to determine the result. This “eager” evaluation should be contrasted with the “lazy” evaluation of `\bool_lazy_... functions`.

T_EXhackers note: The eager evaluation of boolean expressions is unfortunately necessary in T_EX. Indeed, a lazy parser can get confused if `&&` or `||` or parentheses appear as (unbraced)

arguments of some predicates. For instance, the innocuous-looking expression below would break (in a lazy parser) if `#1` were a closing parenthesis and `\l_tmpa_bool` were `true`.

```
( \l_tmpa_bool || \token_if_eq_meaning_p:NN X #1 )
```

Minimal (lazy) evaluation can be obtained using the conditionals `\bool_lazy_all:nTF`, `\bool_lazy_and:nnTF`, `\bool_lazy_any:nTF`, or `\bool_lazy_or:nnTF`, which only evaluate their boolean expression arguments when they are needed to determine the resulting truth value. For example, when evaluating the boolean expression

```
\bool_lazy_and_p:nn
{
  \bool_lazy_any_p:n
  {
    { \int_compare_p:n { 2 = 3 } }
    { \int_compare_p:n { 4 <= 4 } }
    { \int_compare_p:n { 1 = \error } } % skipped
  }
}
{ ! \int_compare_p:n { 2 = 4 } }
```

the line marked with `skipped` is not expanded because the result of `\bool_lazy_any_p:n` is known once the second boolean expression is found to be logically `true`. On the other hand, the last line is expanded because its logical value is needed to determine the result of `\bool_lazy_and_p:nn`.

<code>\bool_if_p:n</code> ★ <code>\bool_if:nTF</code> ★	<code>\bool_if_p:n {<boolean expression>}</code> <code>\bool_if:nTF {<boolean expression>} {<true code>} {<false code>}</code>
--	---

Updated: 2017-07-15 Tests the current truth of *<boolean expression>*, and continues expansion based on this result. The *<boolean expression>* should consist of a series of predicates or boolean variables with the logical relationship between these defined using `&&` (“And”), `||` (“Or”), `!` (“Not”) and parentheses. The logical Not applies to the next predicate or group.

<code>\bool_lazy_all_p:n</code> ★ <code>\bool_lazy_all:nTF</code> ★	<code>\bool_lazy_all_p:n { {<boolexpr₁>} {<boolexpr₂>} ... {<boolexpr_N>} }</code> <code>\bool_lazy_all:nTF { {<boolexpr₁>} {<boolexpr₂>} ... {<boolexpr_N>} } {<true code>} {<false code>}</code>
--	---

New: 2015-11-15
Updated: 2017-07-15 Implements the “And” operation on the *<boolean expressions>*, hence is `true` if all of them are `true` and `false` if any of them is `false`. Contrarily to the infix operator `&&`, only the *<boolean expressions>* which are needed to determine the result of `\bool_lazy_all:nTF` are evaluated. See also `\bool_lazy_and:nnTF` when there are only two *<boolean expressions>*.

<code>\bool_lazy_and_p:nn</code> ★ <code>\bool_lazy_and:nnTF</code> ★	<code>\bool_lazy_and_p:nn {<boolexpr₁>} {<boolexpr₂>}</code> <code>\bool_lazy_and:nnTF {<boolexpr₁>} {<boolexpr₂>} {<true code>} {<false code>}</code>
--	---

New: 2015-11-15
Updated: 2017-07-15 Implements the “And” operation between two boolean expressions, hence is `true` if both are `true`. Contrarily to the infix operator `&&`, the *<boolexpr₂>* is only evaluated if it is needed to determine the result of `\bool_lazy_and:nnTF`. See also `\bool_lazy_all:nTF` when there are more than two *<boolean expressions>*.

<code>\bool_lazy_any_p:n</code> *	<code>\bool_lazy_any_p:n { {<boolean expr₁>} {<boolean expr₂>} ... {<boolean expr_N>} }</code>
<code>\bool_lazy_any:nTF</code> *	<code>\bool_lazy_any:nTF { {<boolean expr₁>} {<boolean expr₂>} ... {<boolean expr_N>} } {<true code>} {<false code>}</code>
New: 2015-11-15	
Updated: 2017-07-15	Implements the “Or” operation on the <i><boolean expressions></i> , hence is true if any of them is true and false if all of them are false . Contrarily to the infix operator <code> </code> , only the <i><boolean expressions></i> which are needed to determine the result of <code>\bool_lazy_any:nTF</code> are evaluated. See also <code>\bool_lazy_or:nnTF</code> when there are only two <i><boolean expressions></i> .

<code>\bool_lazy_or_p:nn</code> *	<code>\bool_lazy_or_p:nn {<boolean expr₁>} {<boolean expr₂>}</code>
<code>\bool_lazy_or:nnTF</code> *	<code>\bool_lazy_or:nnTF {<boolean expr₁>} {<boolean expr₂>} {<true code>} {<false code>}</code>
New: 2015-11-15	
Updated: 2017-07-15	Implements the “Or” operation between two boolean expressions, hence is true if either one is true . Contrarily to the infix operator <code> </code> , the <i><boolean expr₂></i> is only evaluated if it is needed to determine the result of <code>\bool_lazy_or:nnTF</code> . See also <code>\bool_lazy_any:nTF</code> when there are more than two <i><boolean expressions></i> .

<code>\bool_not_p:n</code> *	<code>\bool_not_p:n {<boolean expression>}</code>
Updated: 2017-07-15	Function version of <code>!(<boolean expression>)</code> within a boolean expression.

<code>\bool_xor_p:nn</code> *	<code>\bool_xor_p:nn {<boolean expr₁>} {<boolean expr₂>}</code>
<code>\bool_xor:nnTF</code> *	<code>\bool_xor:nnTF {<boolean expr₁>} {<boolean expr₂>} {<true code>} {<false code>}</code>
New: 2018-05-09	Implements an “exclusive or” operation between two boolean expressions. There is no infix operation for this logical operation.

4 Logical loops

Loops using either boolean expressions or stored boolean values.

<code>\bool_do_until:Nn</code> ☆	<code>\bool_do_until:Nn <boolean> {<code>}</code>
<code>\bool_do_until:cn</code> ☆	
Updated: 2017-07-15	Places the <i><code></i> in the input stream for T _E X to process, and then checks the logical value of the <i><boolean></i> . If it is false then the <i><code></i> is inserted into the input stream again and the process loops until the <i><boolean></i> is true .

<code>\bool_do_while:Nn</code> ☆	<code>\bool_do_while:Nn <boolean> {<code>}</code>
<code>\bool_do_while:cn</code> ☆	
Updated: 2017-07-15	Places the <i><code></i> in the input stream for T _E X to process, and then checks the logical value of the <i><boolean></i> . If it is true then the <i><code></i> is inserted into the input stream again and the process loops until the <i><boolean></i> is false .

<code>\bool_until_do:Nn</code> ☆	<code>\bool_until_do:Nn <boolean> {<code>}</code>
<code>\bool_until_do:cn</code> ☆	
Updated: 2017-07-15	This function firsts checks the logical value of the <i><boolean></i> . If it is false the <i><code></i> is placed in the input stream and expanded. After the completion of the <i><code></i> the truth of the <i><boolean></i> is re-evaluated. The process then loops until the <i><boolean></i> is true .

<code>\bool_while_do:Nn</code> ☆	<code>\bool_while_do:Nn <boolean> {<code>}</code>
<code>\bool_while_do:cn</code> ☆	
Updated: 2017-07-15	This function firsts checks the logical value of the <i><boolean></i> . If it is true the <i><code></i> is placed in the input stream and expanded. After the completion of the <i><code></i> the truth of the <i><boolean></i> is re-evaluated. The process then loops until the <i><boolean></i> is false .

<hr/> <code>\bool_do_until:nn</code> ☆ <hr/>	<code>\bool_do_until:nn {<boolean expression>} {<code>}</code>
Updated: 2017-07-15	Places the <i><code></i> in the input stream for T _E X to process, and then checks the logical value of the <i><boolean expression></i> as described for <code>\bool_if:nTF</code> . If it is false then the <i><code></i> is inserted into the input stream again and the process loops until the <i><boolean expression></i> evaluates to true .
<hr/> <code>\bool_do_while:nn</code> ☆ <hr/>	<code>\bool_do_while:nn {<boolean expression>} {<code>}</code>
Updated: 2017-07-15	Places the <i><code></i> in the input stream for T _E X to process, and then checks the logical value of the <i><boolean expression></i> as described for <code>\bool_if:nTF</code> . If it is true then the <i><code></i> is inserted into the input stream again and the process loops until the <i><boolean expression></i> evaluates to false .
<hr/> <code>\bool_until_do:nn</code> ☆ <hr/>	<code>\bool_until_do:nn {<boolean expression>} {<code>}</code>
Updated: 2017-07-15	This function firsts checks the logical value of the <i><boolean expression></i> (as described for <code>\bool_if:nTF</code>). If it is false the <i><code></i> is placed in the input stream and expanded. After the completion of the <i><code></i> the truth of the <i><boolean expression></i> is re-evaluated. The process then loops until the <i><boolean expression></i> is true .
<hr/> <code>\bool_while_do:nn</code> ☆ <hr/>	<code>\bool_while_do:nn {<boolean expression>} {<code>}</code>
Updated: 2017-07-15	This function firsts checks the logical value of the <i><boolean expression></i> (as described for <code>\bool_if:nTF</code>). If it is true the <i><code></i> is placed in the input stream and expanded. After the completion of the <i><code></i> the truth of the <i><boolean expression></i> is re-evaluated. The process then loops until the <i><boolean expression></i> is false .

5 Producing multiple copies

<hr/> <code>\prg_replicate:nn</code> ★ <hr/>	<code>\prg_replicate:nn {<integer expression>} {<tokens>}</code>
Updated: 2011-07-04	Evaluates the <i><integer expression></i> (which should be zero or positive) and creates the resulting number of copies of the <i><tokens></i> . The function is both expandable and safe for nesting. It yields its result after two expansion steps.

6 Detecting T_EX's mode

<hr/> <code>\mode_if_horizontal_p:</code> ★ <code>\mode_if_horizontal:TF</code> ★ <hr/>	<code>\mode_if_horizontal_p:</code> <code>\mode_if_horizontal:TF {<true code>} {<false code>}</code>
	Detects if T _E X is currently in horizontal mode.
<hr/> <code>\mode_if_inner_p:</code> ★ <code>\mode_if_inner:TF</code> ★ <hr/>	<code>\mode_if_inner_p:</code> <code>\mode_if_inner:TF {<true code>} {<false code>}</code>
	Detects if T _E X is currently in inner mode.
<hr/> <code>\mode_if_math_p:</code> ★ <code>\mode_if_math:TF</code> ★ <hr/>	<code>\mode_if_math:TF {<true code>} {<false code>}</code>
Updated: 2011-09-05	Detects if T _E X is currently in maths mode.

<code>\mode_if_vertical_p: *</code>	<code>\mode_if_vertical_p:</code>
<code>\mode_if_vertical:TF *</code>	<code>\mode_if_vertical:TF {\langle true code \rangle} {\langle false code \rangle}</code>

Detects if T_EX is currently in vertical mode.

7 Primitive conditionals

<code>\if_predicate:w *</code>	<code>\if_predicate:w \langle predicate \rangle \langle true code \rangle \else: \langle false code \rangle \fi:</code>
--------------------------------	---

This function takes a predicate function and branches according to the result. (In practice this function would also accept a single boolean variable in place of the $\langle predicate \rangle$ but to make the coding clearer this should be done through `\if_bool:N`.)

<code>\if_bool:N *</code>	<code>\if_bool:N \langle boolean \rangle \langle true code \rangle \else: \langle false code \rangle \fi:</code>
---------------------------	--

This function takes a boolean variable and branches according to the result.

8 Nestable recursions and mappings

There are a number of places where recursion or mapping constructs are used in `expl3`. At a low-level, these typically require insertion of tokens at the end of the content to allow “clean up”. To support such mappings in a nestable form, the following functions are provided.

<code>\prg_break_point:Nn *</code>	<code>\prg_break_point:Nn \langle type \rangle_map_break: {\langle code \rangle}</code>
------------------------------------	---

New: 2018-03-26

Used to mark the end of a recursion or mapping: the functions `\langle type \rangle_map_break:` and `\langle type \rangle_map_break:n` use this to break out of the loop (see `\prg_map_break:Nn` for how to set these up). After the loop ends, the $\langle code \rangle$ is inserted into the input stream. This occurs even if the break functions are *not* applied: `\prg_break_point:Nn` is functionally-equivalent in these cases to `\use_ii:nn`.

<code>\prg_map_break:Nn *</code>	<code>\prg_map_break:Nn \langle type \rangle_map_break: {\langle user code \rangle}</code>
----------------------------------	--

New: 2018-03-26

`...`
`\prg_break_point:Nn \langle type \rangle_map_break: {\langle ending code \rangle}`

Breaks a recursion in mapping contexts, inserting in the input stream the $\langle user code \rangle$ after the $\langle ending code \rangle$ for the loop. The function breaks loops, inserting their $\langle ending code \rangle$, until reaching a loop with the same $\langle type \rangle$ as its first argument. This `\langle type \rangle_map_break:` argument must be defined; it is simply used as a recognizable marker for the $\langle type \rangle$.

For types with mappings defined in the kernel, `\langle type \rangle_map_break:` and `\langle type \rangle_map_break:n` are defined as `\prg_map_break:Nn \langle type \rangle_map_break: {}` and the same with `{}` omitted.

8.1 Simple mappings

In addition to the more complex mappings above, non-nestable mappings are used in a number of locations and support is provided for these.

<code>\prg_break_point:</code> *	This copy of <code>\prg_do_nothing:</code> is used to mark the end of a fast short-term recursion:
New: 2018-03-27	the function <code>\prg_break:n</code> uses this to break out of the loop.

<code>\prg_break:</code> *	<code>\prg_break:n {<code>} ... \prg_break_point:</code>
<code>\prg_break:n</code> *	Breaks a recursion which has no <i><ending code></i> and which is not a user-breakable mapping
New: 2018-03-27	(see for instance <code>\prop_get:Nn</code>), and inserts the <i><code></i> in the input stream.

9 Internal programming functions

<code>\group_align_safe_begin:</code> *	<code>\group_align_safe_begin:</code>
<code>\group_align_safe_end:</code> *	...
Updated: 2011-08-11	<code>\group_align_safe_end:</code>

These functions are used to enclose material in a T_EX alignment environment within a specially-constructed group. This group is designed in such a way that it does not add brace groups to the output but does act as a group for the `&` token inside `\halign`. This is necessary to allow grabbing of tokens for testing purposes, as T_EX uses group level to determine the effect of alignment tokens. Without the special grouping, the use of a function such as `\peek_after:Nw` would result in a forbidden comparison of the internal `\endtemplate` token, yielding a fatal error. Each `\group_align_safe_begin:` must be matched by a `\group_align_safe_end:`, although this does not have to occur within the same function.

Part XIV

The l3sys package: System/runtime functions

1 The name of the job

`\c_sys_jobname_str`

New: 2015-09-19
Updated: 2019-10-27

Constant that gets the “job name” assigned when T_EX starts.

T_EXhackers note: This copies the contents of the primitive `\jobname`. For technical reasons, the string here is not of the same internal form as other, but may be manipulated using normal string functions.

2 Date and time

`\c_sys_minute_int`
`\c_sys_hour_int`
`\c_sys_day_int`
`\c_sys_month_int`
`\c_sys_year_int`

New: 2015-09-22

The date and time at which the current job was started: these are all reported as integers.

T_EXhackers note: Whilst the underlying primitives can be altered by the user, this interface to the time and date is intended to be the “real” values.

3 Engine

`\sys_if_engine luatex_p:` \star
`\sys_if_engine luatex:` *TF* \star
`\sys_if_engine pdftex_p:` \star
`\sys_if_engine pdftex:` *TF* \star
`\sys_if_engine ptex_p:` \star
`\sys_if_engine ptex:` *TF* \star
`\sys_if_engine uptex_p:` \star
`\sys_if_engine uptex:` *TF* \star
`\sys_if_engine xetex_p:` \star
`\sys_if_engine xetex:` *TF* \star

New: 2015-09-07

`\sys_if_engine pdftex:TF` $\{(true\ code)\} \{(false\ code)\}$

Conditionals which allow engine-specific code to be used. The names follow naturally from those of the engine binaries: note that the (u)ptex tests are for ε -pT_EX and ε -upT_EX as expl3 requires the ε -T_EX extensions. Each conditional is true for *exactly one* supported engine. In particular, `\sys_if_engine ptex_p:` is true for ε -pT_EX but false for ε -upT_EX.

`\c_sys_engine_str`

New: 2015-09-19

The current engine given as a lower case string: one of `luatex`, `pdftex`, `ptex`, `uptex` or `xetex`.

`\c_sys_engine_exec_str`

New: 2020-08-20

The name of the standard executable for the current T_EX engine given as a lower case string: one of `luatex`, `luahtex`, `pdftex`, `eptex`, `euptex` or `xetex`.

`\c_sys_engine_format_str`

New: 2020-08-20

The name of the preloaded format for the current T_EX run given as a lower case string: one of `lualatex` (or `dvilualatex`), `pdflatex` (or `latex`), `platex`, `uplatex` or `xelatex` for L^AT_EX, similar names for plain T_EX (except pdfT_EX in DVI mode yields `etex`), and `cont-en` for ConT_EXt (i.e. the `\fmtname`).

4 Output format

`\sys_if_output_dvi_p: *`
`\sys_if_output_dvi:TF *`
`\sys_if_output_pdf_p: *`
`\sys_if_output_pdf:TF *`

New: 2015-09-19

`\sys_if_output_dvi:TF` `{\true code}` `{\false code}`

Conditionals which give the current output mode the T_EX run is operating in. This is always one of two outcomes, DVI mode or PDF mode. The two sets of conditionals are thus complementary and are both provided to allow the programmer to emphasise the most appropriate case.

`\c_sys_output_str`

New: 2015-09-19

The current output mode given as a lower case string: one of `dvi` or `pdf`.

5 Platform

`\sys_if_platform_unix_p: *` `\sys_if_platform_unix:TF` `{\true code}` `{\false code}`
`\sys_if_platform_unix:TF *`
`\sys_if_platform_windows_p: *`
`\sys_if_platform_windows:TF *`

New: 2018-07-27

Conditionals which allow platform-specific code to be used. The names follow the Lua `os.type()` function, *i.e.* all Unix-like systems are `unix` (including Linux and MacOS).

`\c_sys_platform_str`

New: 2018-07-27

The current platform given as a lower case string: one of `unix`, `windows` or `unknown`.

6 Random numbers

`\sys_rand_seed: *`

New: 2017-05-27

`\sys_rand_seed:`

Expands to the current value of the engine's random seed, a non-negative integer. In engines without random number support this expands to 0.

\sys_gset_rand_seed:n

New: 2017-05-27

\sys_gset_rand_seed:n $\{ \langle \textit{intexpr} \rangle \}$

Globally sets the seed for the engine's pseudo-random number generator to the $\langle \textit{integer expression} \rangle$. This random seed affects all $\backslash \dots _ \textit{rand}$ functions (such as $\backslash \textit{int_rand:nn}$ or $\backslash \textit{clist_rand_item:n}$) as well as other packages relying on the engine's random number generator. In engines without random number support this produces an error.

TeXhackers note: While a 32-bit (signed) integer can be given as a seed, only the absolute value is used and any number beyond 2^{28} is divided by an appropriate power of 2. We recommend using an integer in $[0, 2^{28} - 1]$.

7 Access to the shell

\sys_get_shell:nnN**\sys_get_shell:nnNTF**

New: 2019-09-20

\sys_get_shell:nnN $\{ \langle \textit{shell command} \rangle \}$ $\{ \langle \textit{setup} \rangle \}$ $\langle \textit{tl var} \rangle$ **\sys_get_shell:nnNTF** $\{ \langle \textit{shell command} \rangle \}$ $\{ \langle \textit{setup} \rangle \}$ $\langle \textit{tl var} \rangle$ $\{ \langle \textit{true code} \rangle \}$ $\{ \langle \textit{false code} \rangle \}$

Defines $\langle \textit{tl} \rangle$ to the text returned by the $\langle \textit{shell command} \rangle$. The $\langle \textit{shell command} \rangle$ is converted to a string using $\backslash \textit{tl_to_str:n}$. Category codes may need to be set appropriately via the $\langle \textit{setup} \rangle$ argument, which is run just before running the $\langle \textit{shell command} \rangle$ (in a group). If shell escape is disabled, the $\langle \textit{tl var} \rangle$ will be set to $\backslash \textit{q_no_value}$ in the non-branching version. Note that quote characters (") *cannot* be used inside the $\langle \textit{shell command} \rangle$. The $\backslash \textit{sys_get_shell:nnNTF}$ conditional returns **true** if the shell is available and no quote is detected, and **false** otherwise.

\c_sys_shell_escape_int

New: 2017-05-27

This variable exposes the internal triple of the shell escape status. The possible values are

0 Shell escape is disabled

1 Unrestricted shell escape is enabled

2 Restricted shell escape is enabled

\sys_if_shell_p: ***\sys_if_shell:TF ***

New: 2017-05-27

\sys_if_shell_p:**\sys_if_shell:TF** $\{ \langle \textit{true code} \rangle \}$ $\{ \langle \textit{false code} \rangle \}$

Performs a check for whether shell escape is enabled. This returns true if either of restricted or unrestricted shell escape is enabled.

\sys_if_shell_unrestricted_p: ***\sys_if_shell_unrestricted:TF ***

New: 2017-05-27

\sys_if_shell_unrestricted_p:**\sys_if_shell_unrestricted:TF** $\{ \langle \textit{true code} \rangle \}$ $\{ \langle \textit{false code} \rangle \}$

Performs a check for whether *unrestricted* shell escape is enabled.

<code>\sys_if_shell_restricted_p: *</code>	<code>\sys_if_shell_restricted_p:</code>
<code>\sys_if_shell_restricted:TF *</code>	<code>\sys_if_shell_restricted:TF {\true code} {\false code}</code>

New: 2017-05-27

Performs a check for whether *restricted* shell escape is enabled. This returns false if unrestricted shell escape is enabled. Unrestricted shell escape is not considered a superset of restricted shell escape in this case. To find whether any shell escape is enabled use `\sys_if_shell:.`

<code>\sys_shell_now:n</code>	<code>\sys_shell_now:n {\tokens}</code>
<code>\sys_shell_now:x</code>	

New: 2017-05-27

Execute `\tokens` through shell escape immediately.

<code>\sys_shell_shipout:n</code>	<code>\sys_shell_shipout:n {\tokens}</code>
<code>\sys_shell_shipout:x</code>	

New: 2017-05-27

Execute `\tokens` through shell escape at shipout.

8 Loading configuration data

<code>\sys_load_backend:n</code>	<code>\sys_load_backend:n {\backend}</code>
----------------------------------	---

New: 2019-09-12

Loads the additional configuration file needed for backend support. If the `\backend` is empty, the standard backend for the engine in use will be loaded. This command may only be used once.

<code>\c_sys_backend_str</code>	Set to the name of the backend in use by <code>\sys_load_backend:n</code> when issued.
---------------------------------	--

<code>\sys_load_debug:</code>	<code>\sys_load_debug:</code>
<code>\sys_load_deprecation:</code>	<code>\sys_load_deprecation:</code>

New: 2019-09-12

Load the additional configuration files for debugging support and rolling back deprecations, respectively.

8.1 Final settings

<code>\sys_finalise:</code>	<code>\sys_finalise:</code>
-----------------------------	-----------------------------

New: 2019-10-06

Finalises all system-dependent functionality: required before loading a backend.

Part XV

The `l3clist` package

Comma separated lists

Comma lists contain ordered data where items can be added to the left or right end of the list. This data type allows basic list manipulations such as adding/removing items, applying a function to every item, removing duplicate items, extracting a given item, using the comma list with specified separators, and so on. Sequences (defined in `l3seq`) are safer, faster, and provide more features, so they should often be preferred to comma lists. Comma lists are mostly useful when interfacing with L^AT_EX 2_ε or other code that expects or provides comma list data.

Several items can be added at once. To ease input of comma lists from data provided by a user outside an `\ExplSyntaxOn ... \ExplSyntaxOff` block, spaces are removed from both sides of each comma-delimited argument upon input. Blank arguments are ignored, to allow for trailing commas or repeated commas (which may otherwise arise when concatenating comma lists “by hand”). In addition, a set of braces is removed if the result of space-trimming is braced: this allows the storage of any item in a comma list. For instance,

```
\clist_new:N \l_my_clist
\clist_put_left:Nn \l_my_clist { ~a~ , ~{b}~ , c~\d }
\clist_put_right:Nn \l_my_clist { ~{e}~ , , {{f}} , }
```

results in `\l_my_clist` containing `a,b,c~\d,{e~},{f}}` namely the five items `a`, `b`, `c~\d`, `e~` and `{f}`. Comma lists normally do not contain empty items so the following gives an empty comma list:

```
\clist_clear_new:N \l_my_clist
\clist_put_right:Nn \l_my_clist { , ~ , , }
\clist_if_empty:NTF \l_my_clist { true } { false }
```

and it leaves `true` in the input stream. To include an “unsafe” item (empty, or one that contains a comma, or starts or ends with a space, or is a single brace group), surround it with braces.

Almost all operations on comma lists are noticeably slower than those on sequences so converting the data to sequences using `\seq_set_from_clist:Nn` (see `l3seq`) may be advisable if speed is important. The exception is that `\clist_if_in:NnTF` and `\clist_remove_duplicates:N` may be faster than their sequence analogues for large lists. However, these functions work slowly for “unsafe” items that must be braced, and may produce errors when their argument contains `{`, `}` or `#` (assuming the usual T_EX category codes apply). The sequence data type should thus certainly be preferred to comma lists to store such items.

1 Creating and initialising comma lists

<code>\clist_new:N</code>	<code>\clist_new:N <comma list></code>
<code>\clist_new:c</code>	

Creates a new *<comma list>* or raises an error if the name is already taken. The declaration is global. The *<comma list>* initially contains no items.

<hr/> <code>\clist_const:Nn</code> <code>\clist_const:(Nx cn cx)</code> <hr/> New: 2014-07-05	<code>\clist_const:Nn <clist var> {<comma list>}</code> Creates a new constant <code><clist var></code> or raises an error if the name is already taken. The value of the <code><clist var></code> is set globally to the <code><comma list></code> .
<hr/> <code>\clist_clear:N</code> <code>\clist_clear:c</code> <code>\clist_gclear:N</code> <code>\clist_gclear:c</code> <hr/>	<code>\clist_clear:N <comma list></code> Clears all items from the <code><comma list></code> .
<hr/> <code>\clist_clear_new:N</code> <code>\clist_clear_new:c</code> <code>\clist_gclear_new:N</code> <code>\clist_gclear_new:c</code> <hr/>	<code>\clist_clear_new:N <comma list></code> Ensures that the <code><comma list></code> exists globally by applying <code>\clist_new:N</code> if necessary, then applies <code>\clist_(g)clear:N</code> to leave the list empty.
<hr/> <code>\clist_set_eq:NN</code> <code>\clist_set_eq:(cN Nc cc)</code> <code>\clist_gset_eq:NN</code> <code>\clist_gset_eq:(cN Nc cc)</code> <hr/>	<code>\clist_set_eq:NN <comma list_1> <comma list_2></code> Sets the content of <code><comma list_1></code> equal to that of <code><comma list_2></code> .
<hr/> <code>\clist_set_from_seq:NN</code> <code>\clist_set_from_seq:(cN Nc cc)</code> <code>\clist_gset_from_seq:NN</code> <code>\clist_gset_from_seq:(cN Nc cc)</code> <hr/> New: 2014-07-17	<code>\clist_set_from_seq:NN <comma list> <sequence></code> Converts the data in the <code><sequence></code> into a <code><comma list></code> : the original <code><sequence></code> is unchanged. Items which contain either spaces or commas are surrounded by braces.
<hr/> <code>\clist_concat:NNN</code> <code>\clist_concat:ccc</code> <code>\clist_gconcat:NNN</code> <code>\clist_gconcat:ccc</code> <hr/>	<code>\clist_concat:NNN <comma list_1> <comma list_2> <comma list_3></code> Concatenates the content of <code><comma list_2></code> and <code><comma list_3></code> together and saves the result in <code><comma list_1></code> . The items in <code><comma list_2></code> are placed at the left side of the new comma list.
<hr/> <code>\clist_if_exist_p:N *</code> <code>\clist_if_exist_p:c *</code> <code>\clist_if_exist:NTF *</code> <code>\clist_if_exist:cTF *</code> <hr/> New: 2012-03-03	<code>\clist_if_exist_p:N <comma list></code> <code>\clist_if_exist:NTF <comma list> {<true code>} {<false code>}</code> Tests whether the <code><comma list></code> is currently defined. This does not check that the <code><comma list></code> really is a comma list.

2 Adding data to comma lists

<code>\clist_set:Nn</code>	<code>\clist_set:Nn <comma list> {\langle item_1\rangle,\dots,\langle item_n\rangle}</code>
<code>\clist_set:(NV No Nx cn cV co cx)</code>	
<code>\clist_gset:Nn</code>	
<code>\clist_gset:(NV No Nx cn cV co cx)</code>	

New: 2011-09-06

Sets $\langle comma list \rangle$ to contain the $\langle items \rangle$, removing any previous content from the variable. Blank items are omitted, spaces are removed from both sides of each item, then a set of braces is removed if the resulting space-trimmed item is braced. To store some $\langle tokens \rangle$ as a single $\langle item \rangle$ even if the $\langle tokens \rangle$ contain commas or spaces, add a set of braces: `\clist_set:Nn <comma list> { {\langle tokens \rangle} }`.

<code>\clist_put_left:Nn</code>	<code>\clist_put_left:Nn <comma list> {\langle item_1\rangle,\dots,\langle item_n\rangle}</code>
<code>\clist_put_left:(NV No Nx cn cV co cx)</code>	
<code>\clist_gput_left:Nn</code>	
<code>\clist_gput_left:(NV No Nx cn cV co cx)</code>	

Updated: 2011-09-05

Appends the $\langle items \rangle$ to the left of the $\langle comma list \rangle$. Blank items are omitted, spaces are removed from both sides of each item, then a set of braces is removed if the resulting space-trimmed item is braced. To append some $\langle tokens \rangle$ as a single $\langle item \rangle$ even if the $\langle tokens \rangle$ contain commas or spaces, add a set of braces: `\clist_put_left:Nn <comma list> { {\langle tokens \rangle} }`.

<code>\clist_put_right:Nn</code>	<code>\clist_put_right:Nn <comma list> {\langle item_1\rangle,\dots,\langle item_n\rangle}</code>
<code>\clist_put_right:(NV No Nx cn cV co cx)</code>	
<code>\clist_gput_right:Nn</code>	
<code>\clist_gput_right:(NV No Nx cn cV co cx)</code>	

Updated: 2011-09-05

Appends the $\langle items \rangle$ to the right of the $\langle comma list \rangle$. Blank items are omitted, spaces are removed from both sides of each item, then a set of braces is removed if the resulting space-trimmed item is braced. To append some $\langle tokens \rangle$ as a single $\langle item \rangle$ even if the $\langle tokens \rangle$ contain commas or spaces, add a set of braces: `\clist_put_right:Nn <comma list> { {\langle tokens \rangle} }`.

3 Modifying comma lists

While comma lists are normally used as ordered lists, it may be necessary to modify the content. The functions here may be used to update comma lists, while retaining the order of the unaffected entries.

<code>\clist_remove_duplicates:N</code>	<code>\clist_remove_duplicates:N <comma list></code>
<code>\clist_remove_duplicates:c</code>	
<code>\clist_gremove_duplicates:N</code>	
<code>\clist_gremove_duplicates:c</code>	

Removes duplicate items from the $\langle comma list \rangle$, leaving the left most copy of each item in the $\langle comma list \rangle$. The $\langle item \rangle$ comparison takes place on a token basis, as for `\tl_if_eq:nn(TF)`.

TeXhackers note: This function iterates through every item in the $\langle comma list \rangle$ and does a comparison with the $\langle items \rangle$ already checked. It is therefore relatively slow with large comma lists. Furthermore, it may fail if any of the items in the $\langle comma list \rangle$ contains `{`, `}`, or `#` (assuming the usual TeX category codes apply).

<code>\clist_remove_all:Nn</code>	<code>\clist_remove_all:Nn <comma list> {<item>}</code>
<code>\clist_remove_all:cn</code>	
<code>\clist_gremove_all:Nn</code>	
<code>\clist_gremove_all:cn</code>	

Removes every occurrence of $\langle item \rangle$ from the $\langle comma list \rangle$. The $\langle item \rangle$ comparison takes place on a token basis, as for `\tl_if_eq:nn(TF)`.

Updated: 2011-09-06

TeXhackers note: The function may fail if the $\langle item \rangle$ contains `{`, `}`, or `#` (assuming the usual TeX category codes apply).

<code>\clist_reverse:N</code>	<code>\clist_reverse:N <comma list></code>
<code>\clist_reverse:c</code>	
<code>\clist_greverse:N</code>	
<code>\clist_greverse:c</code>	

New: 2014-07-18

Reverses the order of items stored in the $\langle comma list \rangle$.

<code>\clist_reverse:n</code>	<code>\clist_reverse:n {<comma list>}</code>
-------------------------------	--

New: 2014-07-18

Leaves the items in the $\langle comma list \rangle$ in the input stream in reverse order. Contrarily to other what is done for other n-type $\langle comma list \rangle$ arguments, braces and spaces are preserved by this process.

TeXhackers note: The result is returned within `\unexpanded`, which means that the comma list does not expand further when appearing in an x-type or e-type argument expansion.

<code>\clist_sort:Nn</code>	<code>\clist_sort:Nn <clist var> {<comparison code>}</code>
<code>\clist_sort:cn</code>	
<code>\clist_gsort:Nn</code>	
<code>\clist_gsort:cn</code>	

New: 2017-02-06

Sorts the items in the $\langle clist var \rangle$ according to the $\langle comparison code \rangle$, and assigns the result to $\langle clist var \rangle$. The details of sorting comparison are described in Section 1.

4 Comma list conditionals

<code>\clist_if_empty_p:N</code> *	<code>\clist_if_empty_p:N</code> $\langle comma list \rangle$
<code>\clist_if_empty_p:c</code> *	<code>\clist_if_empty:N</code> $\langle comma list \rangle$ $\{\langle true code \rangle\}$ $\{\langle false code \rangle\}$
<code>\clist_if_empty:N</code> <u><i>TF</i></u> *	Tests if the $\langle comma list \rangle$ is empty (containing no items).
<code>\clist_if_empty:c</code> <u><i>TF</i></u> *	

<code>\clist_if_empty_p:n</code> *	<code>\clist_if_empty_p:n</code> $\{\langle comma list \rangle\}$
<code>\clist_if_empty:n</code> <u><i>TF</i></u> *	<code>\clist_if_empty:n</code> $\langle comma list \rangle$ $\{\langle true code \rangle\}$ $\{\langle false code \rangle\}$

New: 2014-07-05

Tests if the $\langle comma list \rangle$ is empty (containing no items). The rules for space trimming are as for other n-type comma-list functions, hence the comma list $\{\sim, \sim, \sim\}$ (without outer braces) is empty, while $\{\sim, \{ \}, \}$ (without outer braces) contains one element, which happens to be empty: the comma-list is not empty.

<code>\clist_if_in:N</code> <u><i>nnTF</i></u>	<code>\clist_if_in:N</code> $\langle comma list \rangle$ $\langle item \rangle$ $\{\langle true code \rangle\}$ $\{\langle false code \rangle\}$
<code>\clist_if_in:(NV No cn cV co)</code> <u><i>TF</i></u>	
<code>\clist_if_in:nn</code> <u><i>TF</i></u>	
<code>\clist_if_in:(nV no)</code> <u><i>TF</i></u>	

Updated: 2011-09-06

Tests if the $\langle item \rangle$ is present in the $\langle comma list \rangle$. In the case of an n-type $\langle comma list \rangle$, the usual rules of space trimming and brace stripping apply. Hence,

`\clist_if_in:nnTF { a , {b}~ , {b} , c } { b } {true} {false}`

yields true.

T_EXhackers note: The function may fail if the $\langle item \rangle$ contains $\{$, $\}$, or $\#$ (assuming the usual T_EX category codes apply).

5 Mapping to comma lists

The functions described in this section apply a specified function to each item of a comma list. All mappings are done at the current group level, *i.e.* any local assignments made by the $\langle function \rangle$ or $\langle code \rangle$ discussed below remain in effect after the loop.

When the comma list is given explicitly, as an n-type argument, spaces are trimmed around each item. If the result of trimming spaces is empty, the item is ignored. Otherwise, if the item is surrounded by braces, one set is removed, and the result is passed to the mapped function. Thus, if the comma list that is being mapped is $\{a_{_},_{_}\{\{b\}_{_}\}_{_},_{_}\{ \},_{_}\{c\},_{_}\}$ then the arguments passed to the mapped function are ‘a’, ‘{b}_{_}’, an empty argument, and ‘c’.

When the comma list is given as an N-type argument, spaces have already been trimmed on input, and items are simply stripped of one set of braces if any. This case is more efficient than using n-type comma lists.

<code>\clist_map_function:NN</code> ☆	<code>\clist_map_function:NN</code> $\langle comma list \rangle$ $\langle function \rangle$
<code>\clist_map_function:cN</code> ☆	Applies $\langle function \rangle$ to every $\langle item \rangle$ stored in the $\langle comma list \rangle$. The $\langle function \rangle$ receives one argument for each iteration. The $\langle items \rangle$ are returned from left to right. The function
<code>\clist_map_function:nN</code> ☆	<code>\clist_map_inline:Nn</code> is in general more efficient than <code>\clist_map_function:NN</code> .

Updated: 2012-06-29

```
\clist_map_inline:Nn
\clist_map_inline:cn
\clist_map_inline:nn
```

Updated: 2012-06-29

```
\clist_map_inline:Nn <comma list> {<inline function>}
```

Applies *<inline function>* to every *<item>* stored within the *<comma list>*. The *<inline function>* should consist of code which receives the *<item>* as #1. The *<items>* are returned from left to right.

```
\clist_map_variable:NNn
\clist_map_variable:cNn
\clist_map_variable:nNn
```

Updated: 2012-06-29

```
\clist_map_variable:NNn <comma list> <variable> {<code>}
```

Stores each *<item>* of the *<comma list>* in turn in the (token list) *<variable>* and applies the *<code>*. The *<code>* will usually make use of the *<variable>*, but this is not enforced. The assignments to the *<variable>* are local. Its value after the loop is the last *<item>* in the *<comma list>*, or its original value if there were no *<item>*. The *<items>* are returned from left to right.

```
\clist_map_break: ☆
```

Updated: 2012-06-29

```
\clist_map_break:
```

Used to terminate a `\clist_map_...` function before all entries in the *<comma list>* have been processed. This normally takes place within a conditional statement, for example

```
\clist_map_inline:Nn \l_my_clist
{
  \str_if_eq:nnTF { #1 } { bingo }
  { \clist_map_break: }
  {
    % Do something useful
  }
}
```

Use outside of a `\clist_map_...` scenario leads to low level T_EX errors.

T_EXhackers note: When the mapping is broken, additional tokens may be inserted before further items are taken from the input stream. This depends on the design of the mapping function.

`\clist_map_break:n` ☆

Updated: 2012-06-29

`\clist_map_break:n` {<code>}

Used to terminate a `\clist_map_...` function before all entries in the <comma list> have been processed, inserting the <code> after the mapping has ended. This normally takes place within a conditional statement, for example

```
\clist_map_inline:Nn \l_my_clist
{
  \str_if_eq:nnTF { #1 } { bingo }
  { \clist_map_break:n { <code> } }
  {
    % Do something useful
  }
}
```

Use outside of a `\clist_map_...` scenario leads to low level T_EX errors.

T_EXhackers note: When the mapping is broken, additional tokens may be inserted before the <code> is inserted into the input stream. This depends on the design of the mapping function.

`\clist_count:N` ★

`\clist_count:c` ★

`\clist_count:n` ★

New: 2012-07-13

`\clist_count:N` <comma list>

Leaves the number of items in the <comma list> in the input stream as an <integer denotation>. The total number of items in a <comma list> includes those which are duplicates, *i.e.* every item in a <comma list> is counted.

6 Using the content of comma lists directly

`\clist_use:Nnnn` ★

`\clist_use:cnnn` ★

New: 2013-05-26

`\clist_use:Nnnn` <clist var> {<separator between two>}

{<separator between more than two>} {<separator between final two>}

Places the contents of the <clist var> in the input stream, with the appropriate <separator> between the items. Namely, if the comma list has more than two items, the <separator between more than two> is placed between each pair of items except the last, for which the <separator between final two> is used. If the comma list has exactly two items, then they are placed in the input stream separated by the <separator between two>. If the comma list has a single item, it is placed in the input stream, and a comma list with no items produces no output. An error is raised if the variable does not exist or if it is invalid.

For example,

```
\clist_set:Nn \l_tmpa_clist { a , b , , c , {de} , f }
\clist_use:Nnnn \l_tmpa_clist { ~and~ } { ,~ } { ,~and~ }
```

inserts “a, b, c, de, and f” in the input stream. The first separator argument is not used in this case because the comma list has more than 2 items.

T_EXhackers note: The result is returned within the `\unexpanded` primitive (`\exp_not:n`), which means that the <items> do not expand further when appearing in an x-type argument expansion.

`\clist_use:Nn` ★
`\clist_use:cn` ★

New: 2013-05-26

`\clist_use:Nn` $\langle\textit{clist var}\rangle$ $\{\langle\textit{separator}\rangle\}$

Places the contents of the $\langle\textit{clist var}\rangle$ in the input stream, with the $\langle\textit{separator}\rangle$ between the items. If the comma list has a single item, it is placed in the input stream, and a comma list with no items produces no output. An error is raised if the variable does not exist or if it is invalid.

For example,

```
\clist_set:Nn \l_tmpa_clist { a , b , , c , {de} , f }
\clist_use:Nn \l_tmpa_clist { ~and~ }
```

inserts “a and b and c and de and f” in the input stream.

TeXhackers note: The result is returned within the `\unexpanded` primitive (`\exp_not:n`), which means that the $\langle\textit{items}\rangle$ do not expand further when appearing in an `x`-type argument expansion.

7 Comma lists as stacks

Comma lists can be used as stacks, where data is pushed to and popped from the top of the comma list. (The left of a comma list is the top, for performance reasons.) The stack functions for comma lists are not intended to be mixed with the general ordered data functions detailed in the previous section: a comma list should either be used as an ordered data type or as a stack, but not in both ways.

`\clist_get:NN`
`\clist_get:cN`
`\clist_get:NNTF`
`\clist_get:cNTF`

New: 2012-05-14
Updated: 2019-02-16

`\clist_get:NN` $\langle\textit{comma list}\rangle$ $\langle\textit{token list variable}\rangle$

Stores the left-most item from a $\langle\textit{comma list}\rangle$ in the $\langle\textit{token list variable}\rangle$ without removing it from the $\langle\textit{comma list}\rangle$. The $\langle\textit{token list variable}\rangle$ is assigned locally. In the non-branching version, if the $\langle\textit{comma list}\rangle$ is empty the $\langle\textit{token list variable}\rangle$ is set to the marker value `\q_no_value`.

`\clist_pop:NN`
`\clist_pop:cN`

Updated: 2011-09-06

`\clist_pop:NN` $\langle\textit{comma list}\rangle$ $\langle\textit{token list variable}\rangle$

Pops the left-most item from a $\langle\textit{comma list}\rangle$ into the $\langle\textit{token list variable}\rangle$, *i.e.* removes the item from the comma list and stores it in the $\langle\textit{token list variable}\rangle$. Both of the variables are assigned locally.

`\clist_gpop:NN`
`\clist_gpop:cN`

`\clist_gpop:NN` $\langle\textit{comma list}\rangle$ $\langle\textit{token list variable}\rangle$

Pops the left-most item from a $\langle\textit{comma list}\rangle$ into the $\langle\textit{token list variable}\rangle$, *i.e.* removes the item from the comma list and stores it in the $\langle\textit{token list variable}\rangle$. The $\langle\textit{comma list}\rangle$ is modified globally, while the assignment of the $\langle\textit{token list variable}\rangle$ is local.

`\clist_pop:NNTF`
`\clist_pop:cNTF`

New: 2012-05-14

`\clist_pop:NNTF` $\langle\textit{comma list}\rangle$ $\langle\textit{token list variable}\rangle$ $\{\langle\textit{true code}\rangle\}$ $\{\langle\textit{false code}\rangle\}$

If the $\langle\textit{comma list}\rangle$ is empty, leaves the $\langle\textit{false code}\rangle$ in the input stream. The value of the $\langle\textit{token list variable}\rangle$ is not defined in this case and should not be relied upon. If the $\langle\textit{comma list}\rangle$ is non-empty, pops the top item from the $\langle\textit{comma list}\rangle$ in the $\langle\textit{token list variable}\rangle$, *i.e.* removes the item from the $\langle\textit{comma list}\rangle$. Both the $\langle\textit{comma list}\rangle$ and the $\langle\textit{token list variable}\rangle$ are assigned locally.

<hr/> <code>\clist_gpop:NNTF</code> <hr/>	<code>\clist_gpop:NNTF <comma list> <token list variable> {<true code>} {<false code>}</code>
<code>\clist_gpop:cNTF</code> <hr/>	
New: 2012-05-14	

If the *<comma list>* is empty, leaves the *<false code>* in the input stream. The value of the *<token list variable>* is not defined in this case and should not be relied upon. If the *<comma list>* is non-empty, pops the top item from the *<comma list>* in the *<token list variable>*, *i.e.* removes the item from the *<comma list>*. The *<comma list>* is modified globally, while the *<token list variable>* is assigned locally.

<hr/> <code>\clist_push:Nn</code> <hr/>	<code>\clist_push:Nn <comma list> {<items>}</code>
<code>\clist_push:(NV No Nx cn cV co cx)</code>	
<code>\clist_gpush:Nn</code>	
<code>\clist_gpush:(NV No Nx cn cV co cx)</code> <hr/>	

Adds the *{<items>}* to the top of the *<comma list>*. Spaces are removed from both sides of each item as for any n-type comma list.

8 Using a single item

<hr/> <code>\clist_item:Nn *</code> <hr/>	<code>\clist_item:Nn <comma list> {<integer expression>}</code>
<code>\clist_item:cn *</code>	
<code>\clist_item:nn *</code> <hr/>	
New: 2014-07-17	

Indexing items in the *<comma list>* from 1 at the top (left), this function evaluates the *<integer expression>* and leaves the appropriate item from the comma list in the input stream. If the *<integer expression>* is negative, indexing occurs from the bottom (right) of the comma list. When the *<integer expression>* is larger than the number of items in the *<comma list>* (as calculated by `\clist_count:N`) then the function expands to nothing.

TeXhackers note: The result is returned within the `\unexpanded` primitive (`\exp_not:n`), which means that the *<item>* does not expand further when appearing in an x-type argument expansion.

<hr/> <code>\clist_rand_item:N *</code> <hr/>	<code>\clist_rand_item:N <clist var></code>
<code>\clist_rand_item:c *</code>	<code>\clist_rand_item:n {<comma list>}</code>
<code>\clist_rand_item:n *</code> <hr/>	
New: 2016-12-06	

Selects a pseudo-random item of the *<comma list>*. If the *<comma list>* has no item, the result is empty.

TeXhackers note: The result is returned within the `\unexpanded` primitive (`\exp_not:n`), which means that the *<item>* does not expand further when appearing in an x-type argument expansion.

9 Viewing comma lists

<hr/> <code>\clist_show:N</code> <hr/>	<code>\clist_show:N <comma list></code>
<code>\clist_show:c</code> <hr/>	
Updated: 2015-08-03	

Displays the entries in the *<comma list>* in the terminal.

<hr/> <code>\clist_show:n</code> <hr/>	<code>\clist_show:n {\tokens}</code>
Updated: 2013-08-03	Displays the entries in the comma list in the terminal.
<hr/>	
<code>\clist_log:N</code> <code>\clist_log:c</code> <hr/>	<code>\clist_log:N <comma list></code>
New: 2014-08-22 Updated: 2015-08-03	Writes the entries in the <i><comma list></i> in the log file. See also <code>\clist_show:N</code> which displays the result in the terminal.
<hr/>	
<code>\clist_log:n</code> <hr/>	<code>\clist_log:n {\tokens}</code>
New: 2014-08-22	Writes the entries in the comma list in the log file. See also <code>\clist_show:n</code> which displays the result in the terminal.

10 Constant and scratch comma lists

<hr/> <code>\c_empty_clist</code> <hr/>	Constant that is always empty.
New: 2012-07-02	
<hr/>	
<code>\l_tmpa_clist</code> <code>\l_tmpb_clist</code> <hr/>	Scratch comma lists for local assignment. These are never used by the kernel code, and so are safe for use with any L ^A T _E X3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.
New: 2011-09-06	
<hr/>	
<code>\g_tmpa_clist</code> <code>\g_tmpb_clist</code> <hr/>	Scratch comma lists for global assignment. These are never used by the kernel code, and so are safe for use with any L ^A T _E X3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.
New: 2011-09-06	

Part XVI

The l3token package

Token manipulation

This module deals with tokens. Now this is perhaps not the most precise description so let's try with a better description: When programming in T_EX, it is often desirable to know just what a certain token is: is it a control sequence or something else. Similarly one often needs to know if a control sequence is expandable or not, a macro or a primitive, how many arguments it takes etc. Another thing of great importance (especially when it comes to document commands) is looking ahead in the token stream to see if a certain character is present and maybe even remove it or disregard other tokens while scanning. This module provides functions for both and as such has two primary function categories: `\token_` for anything that deals with tokens and `\peek_` for looking ahead in the token stream.

Most functions we describe here can be used on control sequences, as those are tokens as well.

It is important to distinguish two aspects of a token: its “shape” (for lack of a better word), which affects the matching of delimited arguments and the comparison of token lists containing this token, and its “meaning”, which affects whether the token expands or what operation it performs. One can have tokens of different shapes with the same meaning, but not the converse.

For instance, `\if:w`, `\if_charcode:w`, and `\tex_if:D` are three names for the same internal operation of T_EX, namely the primitive testing the next two characters for equality of their character code. They have the same meaning hence behave identically in many situations. However, T_EX distinguishes them when searching for a delimited argument. Namely, the example function `\show_until_if:w` defined below takes everything until `\if:w` as an argument, despite the presence of other copies of `\if:w` under different names.

```
\cs_new:Npn \show_until_if:w #1 \if:w { \tl_show:n {#1} }
\show_until_if:w \tex_if:D \if_charcode:w \if:w
```

A list of all possible shapes and a list of all possible meanings are given in section 7.

1 Creating character tokens

```
\char_set_active_eq:NN
\char_set_active_eq:Nc
\char_gset_active_eq:NN
\char_gset_active_eq:Nc
```

Updated: 2015-11-12

```
\char_set_active_eq:NN <char> <function>
```

Sets the behaviour of the `<char>` in situations where it is active (category code 13) to be equivalent to that of the `<function>`. The category code of the `<char>` is *unchanged* by this process. The `<function>` may itself be an active character.

```
\char_set_active_eq:nN
\char_set_active_eq:nc
\char_gset_active_eq:nN
\char_gset_active_eq:nc
```

New: 2015-11-12

```
\char_set_active_eq:nN {<integer expression>} <function>
```

Sets the behaviour of the `<char>` which has character code as given by the `<integer expression>` in situations where it is active (category code 13) to be equivalent to that of the `<function>`. The category code of the `<char>` is *unchanged* by this process. The `<function>` may itself be an active character.

<hr/> <code>\char_generate:nn</code> ★ <hr/>	<code>\char_generate:nn</code> { $\langle charcode \rangle$ } { $\langle catcode \rangle$ }
New: 2015-09-09 Updated: 2019-01-16	Generates a character token of the given $\langle charcode \rangle$ and $\langle catcode \rangle$ (both of which may be integer expressions). The $\langle catcode \rangle$ may be one of

- 1 (begin group)
- 2 (end group)
- 3 (math toggle)
- 4 (alignment)
- 6 (parameter)
- 7 (math superscript)
- 8 (math subscript)
- 11 (letter)
- 12 (other)
- 13 (active)

and other values raise an error. The $\langle charcode \rangle$ may be any one valid for the engine in use. Active characters cannot be generated in older versions of X_YTeX.

T_EXhackers note: Exactly two expansions are needed to produce the character.

<hr/> <code>\char_lowercase:N</code> ★ <hr/>	<code>\char_lowercase:N</code> $\langle char \rangle$
<code>\char_uppercase:N</code> ★	Converts the $\langle char \rangle$ to the equivalent case-changed character as detailed by the function name (see <code>\str_foldcase:n</code> and <code>\text_titlecase:n</code> for details of these terms). The case mapping is carried out with no context-dependence (<i>cf.</i> <code>\text_uppercase:n</code> , <i>etc.</i>) The str versions always generate “other” (category code 12) characters, whilst the standard versions generate characters with the category code of the $\langle char \rangle$ (i.e. only the character code changes).
<code>\char_titlecase:N</code> ★	
<code>\char_foldcase:N</code> ★	
<code>\char_str_lowercase:N</code> ★	
<code>\char_str_uppercase:N</code> ★	
<code>\char_str_titlecase:N</code> ★	
<code>\char_str_foldcase:N</code> ★	
New: 2020-01-09	

<hr/> <code>\c_catcode_other_space_tl</code> <hr/>	Token list containing one character with category code 12, (“other”), and character code 32 (space).
New: 2011-09-05	

2 Manipulating and interrogating character tokens

<code>\char_set_catcode_escape:N</code>	<code>\char_set_catcode_letter:N</code> $\langle character \rangle$
<code>\char_set_catcode_group_begin:N</code>	
<code>\char_set_catcode_group_end:N</code>	
<code>\char_set_catcode_math_toggle:N</code>	
<code>\char_set_catcode_alignment:N</code>	
<code>\char_set_catcode_end_line:N</code>	
<code>\char_set_catcode_parameter:N</code>	
<code>\char_set_catcode_math_superscript:N</code>	
<code>\char_set_catcode_math_subscript:N</code>	
<code>\char_set_catcode_ignore:N</code>	
<code>\char_set_catcode_space:N</code>	
<code>\char_set_catcode_letter:N</code>	
<code>\char_set_catcode_other:N</code>	
<code>\char_set_catcode_active:N</code>	
<code>\char_set_catcode_comment:N</code>	
<code>\char_set_catcode_invalid:N</code>	

Updated: 2015-11-11

Sets the category code of the $\langle character \rangle$ to that indicated in the function name. Depending on the current category code of the $\langle token \rangle$ the escape token may also be needed:

`\char_set_catcode_other:N \%`

The assignment is local.

<code>\char_set_catcode_escape:n</code>	<code>\char_set_catcode_letter:n</code> $\{ \langle integer\ expression \rangle \}$
<code>\char_set_catcode_group_begin:n</code>	
<code>\char_set_catcode_group_end:n</code>	
<code>\char_set_catcode_math_toggle:n</code>	
<code>\char_set_catcode_alignment:n</code>	
<code>\char_set_catcode_end_line:n</code>	
<code>\char_set_catcode_parameter:n</code>	
<code>\char_set_catcode_math_superscript:n</code>	
<code>\char_set_catcode_math_subscript:n</code>	
<code>\char_set_catcode_ignore:n</code>	
<code>\char_set_catcode_space:n</code>	
<code>\char_set_catcode_letter:n</code>	
<code>\char_set_catcode_other:n</code>	
<code>\char_set_catcode_active:n</code>	
<code>\char_set_catcode_comment:n</code>	
<code>\char_set_catcode_invalid:n</code>	

Updated: 2015-11-11

Sets the category code of the $\langle character \rangle$ which has character code as given by the $\langle integer\ expression \rangle$. This version can be used to set up characters which cannot otherwise be given (*cf.* the N-type variants). The assignment is local.

<hr/> <code>\char_set_catcode:nn</code> <hr/>	<code>\char_set_catcode:nn {⟨intexpr₁⟩} {⟨intexpr₂⟩}</code>
<hr/> Updated: 2015-11-11 <hr/>	These functions set the category code of the <i>⟨character⟩</i> which has character code as given by the <i>⟨integer expression⟩</i> . The first <i>⟨integer expression⟩</i> is the character code and the second is the category code to apply. The setting applies within the current T _E X group. In general, the symbolic functions <code>\char_set_catcode_⟨type⟩</code> should be preferred, but there are cases where these lower-level functions may be useful.
<hr/> <code>\char_value_catcode:n ★</code> <hr/>	<code>\char_value_catcode:n {⟨integer expression⟩}</code>
	Expands to the current category code of the <i>⟨character⟩</i> with character code given by the <i>⟨integer expression⟩</i> .
<hr/> <code>\char_show_value_catcode:n</code> <hr/>	<code>\char_show_value_catcode:n {⟨integer expression⟩}</code>
	Displays the current category code of the <i>⟨character⟩</i> with character code given by the <i>⟨integer expression⟩</i> on the terminal.
<hr/> <code>\char_set_lccode:nn</code> <hr/>	<code>\char_set_lccode:nn {⟨intexpr₁⟩} {⟨intexpr₂⟩}</code>
<hr/> Updated: 2015-08-06 <hr/>	Sets up the behaviour of the <i>⟨character⟩</i> when found inside <code>\text_lowercase:n</code> , such that <i>⟨character₁⟩</i> will be converted into <i>⟨character₂⟩</i> . The two <i>⟨characters⟩</i> may be specified using an <i>⟨integer expression⟩</i> for the character code concerned. This may include the T _E X ‘ <i>⟨character⟩</i> ’ method for converting a single character into its character code:
	<pre> \char_set_lccode:nn { ‘\A } { ‘a } % Standard behaviour \char_set_lccode:nn { ‘\A } { ‘\A + 32 } \char_set_lccode:nn { 50 } { 60 } </pre>
	The setting applies within the current T _E X group.
<hr/> <code>\char_value_lccode:n ★</code> <hr/>	<code>\char_value_lccode:n {⟨integer expression⟩}</code>
	Expands to the current lower case code of the <i>⟨character⟩</i> with character code given by the <i>⟨integer expression⟩</i> .
<hr/> <code>\char_show_value_lccode:n</code> <hr/>	<code>\char_show_value_lccode:n {⟨integer expression⟩}</code>
	Displays the current lower case code of the <i>⟨character⟩</i> with character code given by the <i>⟨integer expression⟩</i> on the terminal.
<hr/> <code>\char_set_uccode:nn</code> <hr/>	<code>\char_set_uccode:nn {⟨intexpr₁⟩} {⟨intexpr₂⟩}</code>
<hr/> Updated: 2015-08-06 <hr/>	Sets up the behaviour of the <i>⟨character⟩</i> when found inside <code>\text_uppercase:n</code> , such that <i>⟨character₁⟩</i> will be converted into <i>⟨character₂⟩</i> . The two <i>⟨characters⟩</i> may be specified using an <i>⟨integer expression⟩</i> for the character code concerned. This may include the T _E X ‘ <i>⟨character⟩</i> ’ method for converting a single character into its character code:
	<pre> \char_set_uccode:nn { ‘a } { ‘\A } % Standard behaviour \char_set_uccode:nn { ‘\A } { ‘\A - 32 } \char_set_uccode:nn { 60 } { 50 } </pre>
	The setting applies within the current T _E X group.

<hr/> <code>\char_value_uccode:n</code> ★ <hr/>	<code>\char_value_uccode:n {⟨integer expression⟩}</code> Expands to the current upper case code of the <i>⟨character⟩</i> with character code given by the <i>⟨integer expression⟩</i> .
<hr/> <code>\char_show_value_uccode:n</code> <hr/>	<code>\char_show_value_uccode:n {⟨integer expression⟩}</code> Displays the current upper case code of the <i>⟨character⟩</i> with character code given by the <i>⟨integer expression⟩</i> on the terminal.
<hr/> <code>\char_set_mathcode:nn</code> <div>Updated: 2015-08-06</div> <hr/>	<code>\char_set_mathcode:nn {⟨intexpr₁⟩} {⟨intexpr₂⟩}</code> This function sets up the math code of <i>⟨character⟩</i> . The <i>⟨character⟩</i> is specified as an <i>⟨integer expression⟩</i> which will be used as the character code of the relevant character. The setting applies within the current T _E X group.
<hr/> <code>\char_value_mathcode:n</code> ★ <hr/>	<code>\char_value_mathcode:n {⟨integer expression⟩}</code> Expands to the current math code of the <i>⟨character⟩</i> with character code given by the <i>⟨integer expression⟩</i> .
<hr/> <code>\char_show_value_mathcode:n</code> <hr/>	<code>\char_show_value_mathcode:n {⟨integer expression⟩}</code> Displays the current math code of the <i>⟨character⟩</i> with character code given by the <i>⟨integer expression⟩</i> on the terminal.
<hr/> <code>\char_set_sfcode:nn</code> <div>Updated: 2015-08-06</div> <hr/>	<code>\char_set_sfcode:nn {⟨intexpr₁⟩} {⟨intexpr₂⟩}</code> This function sets up the space factor for the <i>⟨character⟩</i> . The <i>⟨character⟩</i> is specified as an <i>⟨integer expression⟩</i> which will be used as the character code of the relevant character. The setting applies within the current T _E X group.
<hr/> <code>\char_value_sfcode:n</code> ★ <hr/>	<code>\char_value_sfcode:n {⟨integer expression⟩}</code> Expands to the current space factor for the <i>⟨character⟩</i> with character code given by the <i>⟨integer expression⟩</i> .
<hr/> <code>\char_show_value_sfcode:n</code> <hr/>	<code>\char_show_value_sfcode:n {⟨integer expression⟩}</code> Displays the current space factor for the <i>⟨character⟩</i> with character code given by the <i>⟨integer expression⟩</i> on the terminal.
<hr/> <code>\l_char_active_seq</code> <div>New: 2012-01-23 Updated: 2015-11-11</div> <hr/>	Used to track which tokens may require special handling at the document level as they are (or have been at some point) of category <i>⟨active⟩</i> (catcode 13). Each entry in the sequence consists of a single escaped token, for example <code>\~</code> . Active tokens should be added to the sequence when they are defined for general document use.
<hr/> <code>\l_char_special_seq</code> <div>New: 2012-01-23 Updated: 2015-11-11</div> <hr/>	Used to track which tokens will require special handling when working with verbatim-like material at the document level as they are not of categories <i>⟨letter⟩</i> (catcode 11) or <i>⟨other⟩</i> (catcode 12). Each entry in the sequence consists of a single escaped token, for example <code>\</code> for the backslash or <code>{</code> for an opening brace. Escaped tokens should be added to the sequence when they are defined for general document use.

3 Generic tokens

```
\c_group_begin_token
\c_group_end_token
\c_math_toggle_token
\c_alignment_token
\c_parameter_token
\c_math_superscript_token
\c_math_subscript_token
\c_space_token
```

These are implicit tokens which have the category code described by their name. They are used internally for test purposes but are also available to the programmer for other uses.

```
\c_catcode_letter_token
\c_catcode_other_token
```

These are implicit tokens which have the category code described by their name. They are used internally for test purposes and should not be used other than for category code tests.

```
\c_catcode_active_tl
```

A token list containing an active token. This is used internally for test purposes and should not be used other than in appropriately-constructed category code tests.

4 Converting tokens

```
\token_to_meaning:N ★
\token_to_meaning:c ★
```

`\token_to_meaning:N` $\langle token \rangle$

Inserts the current meaning of the $\langle token \rangle$ into the input stream as a series of characters of category code 12 (other). This is the primitive \TeX description of the $\langle token \rangle$, thus for example both functions defined by `\cs_set_nopar:Npn` and token list variables defined using `\tl_new:N` are described as macros.

\TeX hackers note: This is the \TeX primitive `\meaning`. The $\langle token \rangle$ can thus be an explicit space tokens or an explicit begin-group or end-group character token (`{` or `}` when normal \TeX category codes apply) even though these are not valid N-type arguments.

```
\token_to_str:N ★
\token_to_str:c ★
```

`\token_to_str:N` $\langle token \rangle$

Converts the given $\langle token \rangle$ into a series of characters with category code 12 (other). If the $\langle token \rangle$ is a control sequence, this will start with the current escape character with category code 12 (the escape character is part of the $\langle token \rangle$). This function requires only a single expansion.

\TeX hackers note: `\token_to_str:N` is the \TeX primitive `\string` renamed. The $\langle token \rangle$ can thus be an explicit space tokens or an explicit begin-group or end-group character token (`{` or `}` when normal \TeX category codes apply) even though these are not valid N-type arguments.

5 Token conditionals

<code>\token_if_group_begin_p:N</code>	<code>*</code>	<code>\token_if_group_begin_p:N</code>	<code><token></code>
<code>\token_if_group_begin:NTF</code>	<code>*</code>	<code>\token_if_group_begin:NTF</code>	<code><token> {\true code} {\false code}</code>

Tests if `<token>` has the category code of a begin group token (`{` when normal `TEX` category codes are in force). Note that an explicit begin group token cannot be tested in this way, as it is not a valid N-type argument.

<code>\token_if_group_end_p:N</code>	<code>*</code>	<code>\token_if_group_end_p:N</code>	<code><token></code>
<code>\token_if_group_end:NTF</code>	<code>*</code>	<code>\token_if_group_end:NTF</code>	<code><token> {\true code} {\false code}</code>

Tests if `<token>` has the category code of an end group token (`}` when normal `TEX` category codes are in force). Note that an explicit end group token cannot be tested in this way, as it is not a valid N-type argument.

<code>\token_if_math_toggle_p:N</code>	<code>*</code>	<code>\token_if_math_toggle_p:N</code>	<code><token></code>
<code>\token_if_math_toggle:NTF</code>	<code>*</code>	<code>\token_if_math_toggle:NTF</code>	<code><token> {\true code} {\false code}</code>

Tests if `<token>` has the category code of a math shift token (`$` when normal `TEX` category codes are in force).

<code>\token_if_alignment_p:N</code>	<code>*</code>	<code>\token_if_alignment_p:N</code>	<code><token></code>
<code>\token_if_alignment:NTF</code>	<code>*</code>	<code>\token_if_alignment:NTF</code>	<code><token> {\true code} {\false code}</code>

Tests if `<token>` has the category code of an alignment token (`&` when normal `TEX` category codes are in force).

<code>\token_if_parameter_p:N</code>	<code>*</code>	<code>\token_if_parameter_p:N</code>	<code><token></code>
<code>\token_if_parameter:NTF</code>	<code>*</code>	<code>\token_if_parameter:NTF</code>	<code><token> {\true code} {\false code}</code>

Tests if `<token>` has the category code of a macro parameter token (`#` when normal `TEX` category codes are in force).

<code>\token_if_math_superscript_p:N</code>	<code>*</code>	<code>\token_if_math_superscript_p:N</code>	<code><token></code>
<code>\token_if_math_superscript:NTF</code>	<code>*</code>	<code>\token_if_math_superscript:NTF</code>	<code><token> {\true code} {\false code}</code>

Tests if `<token>` has the category code of a superscript token (`^` when normal `TEX` category codes are in force).

<code>\token_if_math_subscript_p:N</code>	<code>*</code>	<code>\token_if_math_subscript_p:N</code>	<code><token></code>
<code>\token_if_math_subscript:NTF</code>	<code>*</code>	<code>\token_if_math_subscript:NTF</code>	<code><token> {\true code} {\false code}</code>

Tests if `<token>` has the category code of a subscript token (`_` when normal `TEX` category codes are in force).

<code>\token_if_space_p:N</code>	<code>*</code>	<code>\token_if_space_p:N</code>	<code><token></code>
<code>\token_if_space:NTF</code>	<code>*</code>	<code>\token_if_space:NTF</code>	<code><token> {\true code} {\false code}</code>

Tests if `<token>` has the category code of a space token. Note that an explicit space token with character code 32 cannot be tested in this way, as it is not a valid N-type argument.

<code>\token_if_letter_p:N</code>	<code>\token_if_letter_p:N</code>	<code>\token</code>
<code>\token_if_letter:NTF</code>	<code>\token_if_letter:NTF</code>	<code>\token</code> <code>{\true code}</code> <code>{\false code}</code>

Tests if `\token` has the category code of a letter token.

<code>\token_if_other_p:N</code>	<code>\token_if_other_p:N</code>	<code>\token</code>
<code>\token_if_other:NTF</code>	<code>\token_if_other:NTF</code>	<code>\token</code> <code>{\true code}</code> <code>{\false code}</code>

Tests if `\token` has the category code of an “other” token.

<code>\token_if_active_p:N</code>	<code>\token_if_active_p:N</code>	<code>\token</code>
<code>\token_if_active:NTF</code>	<code>\token_if_active:NTF</code>	<code>\token</code> <code>{\true code}</code> <code>{\false code}</code>

Tests if `\token` has the category code of an active character.

<code>\token_if_eq_catcode_p:NN</code>	<code>\token_if_eq_catcode_p:NN</code>	<code>\token₁</code>	<code>\token₂</code>
<code>\token_if_eq_catcode:NNTF</code>	<code>\token_if_eq_catcode:NNTF</code>	<code>\token₁</code>	<code>\token₂</code> <code>{\true code}</code> <code>{\false code}</code>

Tests if the two `\tokens` have the same category code.

<code>\token_if_eq_charcode_p:NN</code>	<code>\token_if_eq_charcode_p:NN</code>	<code>\token₁</code>	<code>\token₂</code>
<code>\token_if_eq_charcode:NNTF</code>	<code>\token_if_eq_charcode:NNTF</code>	<code>\token₁</code>	<code>\token₂</code> <code>{\true code}</code> <code>{\false code}</code>

Tests if the two `\tokens` have the same character code.

<code>\token_if_eq_meaning_p:NN</code>	<code>\token_if_eq_meaning_p:NN</code>	<code>\token₁</code>	<code>\token₂</code>
<code>\token_if_eq_meaning:NNTF</code>	<code>\token_if_eq_meaning:NNTF</code>	<code>\token₁</code>	<code>\token₂</code> <code>{\true code}</code> <code>{\false code}</code>

Tests if the two `\tokens` have the same meaning when expanded.

<code>\token_if_macro_p:N</code>	<code>\token_if_macro_p:N</code>	<code>\token</code>
<code>\token_if_macro:NTF</code>	<code>\token_if_macro:NTF</code>	<code>\token</code> <code>{\true code}</code> <code>{\false code}</code>

Updated: 2011-05-23 Tests if the `\token` is a \TeX macro.

<code>\token_if_cs_p:N</code>	<code>\token_if_cs_p:N</code>	<code>\token</code>
<code>\token_if_cs:NTF</code>	<code>\token_if_cs:NTF</code>	<code>\token</code> <code>{\true code}</code> <code>{\false code}</code>

Tests if the `\token` is a control sequence.

<code>\token_if_expandable_p:N</code>	<code>\token_if_expandable_p:N</code>	<code>\token</code>
<code>\token_if_expandable:NTF</code>	<code>\token_if_expandable:NTF</code>	<code>\token</code> <code>{\true code}</code> <code>{\false code}</code>

Tests if the `\token` is expandable. This test returns `\false` for an undefined token.

<code>\token_if_long_macro_p:N</code>	<code>\token_if_long_macro_p:N</code>	<code>\token</code>
<code>\token_if_long_macro:NTF</code>	<code>\token_if_long_macro:NTF</code>	<code>\token</code> <code>{\true code}</code> <code>{\false code}</code>

Updated: 2012-01-20 Tests if the `\token` is a long macro.

<code>\token_if_protected_macro_p:N</code>	<code>\token_if_protected_macro_p:N</code>	<code>\token</code>
<code>\token_if_protected_macro:NTF</code>	<code>\token_if_protected_macro:NTF</code>	<code>\token</code> <code>{\true code}</code> <code>{\false code}</code>

Updated: 2012-01-20

Tests if the `\token` is a protected macro: for a macro which is both protected and long this returns `false`.

<code>\token_if_protected_long_macro_p:N</code>	<code>\token_if_protected_long_macro_p:N</code>	<code>\token</code>
<code>\token_if_protected_long_macro:NTF</code>	<code>\token_if_protected_long_macro:NTF</code>	<code>\token</code> <code>{\true code}</code> <code>{\false code}</code>

Updated: 2012-01-20

Tests if the `\token` is a protected long macro.

<code>\token_if_chardef_p:N</code>	<code>\token_if_chardef_p:N</code>	<code>\token</code>
<code>\token_if_chardef:NTF</code>	<code>\token_if_chardef:NTF</code>	<code>\token</code> <code>{\true code}</code> <code>{\false code}</code>

Updated: 2012-01-20

Tests if the `\token` is defined to be a chardef.

T_EXhackers note: Booleans, boxes and small integer constants are implemented as `\chardefs`.

<code>\token_if_mathchardef_p:N</code>	<code>\token_if_mathchardef_p:N</code>	<code>\token</code>
<code>\token_if_mathchardef:NTF</code>	<code>\token_if_mathchardef:NTF</code>	<code>\token</code> <code>{\true code}</code> <code>{\false code}</code>

Updated: 2012-01-20

Tests if the `\token` is defined to be a mathchardef.

<code>\token_if_font_selection_p:N</code>	<code>\token_if_font_selection_p:N</code>	<code>\token</code>
<code>\token_if_font_selection:NTF</code>	<code>\token_if_font_selection:NTF</code>	<code>\token</code> <code>{\true code}</code> <code>{\false code}</code>

New: 2020-10-27

Tests if the `\token` is defined to be a font selection command.

<code>\token_if_dim_register_p:N</code>	<code>\token_if_dim_register_p:N</code>	<code>\token</code>
<code>\token_if_dim_register:NTF</code>	<code>\token_if_dim_register:NTF</code>	<code>\token</code> <code>{\true code}</code> <code>{\false code}</code>

Updated: 2012-01-20

Tests if the `\token` is defined to be a dimension register.

<code>\token_if_int_register_p:N</code>	<code>\token_if_int_register_p:N</code>	<code>\token</code>
<code>\token_if_int_register:NTF</code>	<code>\token_if_int_register:NTF</code>	<code>\token</code> <code>{\true code}</code> <code>{\false code}</code>

Updated: 2012-01-20

Tests if the `\token` is defined to be a integer register.

T_EXhackers note: Constant integers may be implemented as integer registers, `\chardefs`, or `\mathchardefs` depending on their value.

<code>\token_if_muskip_register_p:N</code>	<code>\token_if_muskip_register_p:N</code>	<code>\token</code>
<code>\token_if_muskip_register:NTF</code>	<code>\token_if_muskip_register:NTF</code>	<code>\token</code> <code>{\true code}</code> <code>{\false code}</code>

New: 2012-02-15

Tests if the `\token` is defined to be a muskip register.

<code>\token_if_skip_register_p:N</code>	<code>*</code>	<code>\token_if_skip_register_p:N</code>	$\langle token \rangle$
<code>\token_if_skip_register:NTF</code>	<code>*</code>	<code>\token_if_skip_register:NTF</code>	$\langle token \rangle$ $\{\langle true\ code \rangle\}$ $\{\langle false\ code \rangle\}$

Updated: 2012-01-20

Tests if the $\langle token \rangle$ is defined to be a skip register.

<code>\token_if_toks_register_p:N</code>	<code>*</code>	<code>\token_if_toks_register_p:N</code>	$\langle token \rangle$
<code>\token_if_toks_register:NTF</code>	<code>*</code>	<code>\token_if_toks_register:NTF</code>	$\langle token \rangle$ $\{\langle true\ code \rangle\}$ $\{\langle false\ code \rangle\}$

Updated: 2012-01-20

Tests if the $\langle token \rangle$ is defined to be a toks register (not used by L^AT_EX3).

<code>\token_if_primitive_p:N</code>	<code>*</code>	<code>\token_if_primitive_p:N</code>	$\langle token \rangle$
<code>\token_if_primitive:NTF</code>	<code>*</code>	<code>\token_if_primitive:NTF</code>	$\langle token \rangle$ $\{\langle true\ code \rangle\}$ $\{\langle false\ code \rangle\}$

Updated: 2020-09-11

Tests if the $\langle token \rangle$ is an engine primitive. In LuaT_EX this includes primitive-like commands defined using `\token.set_lua`.

<code>\token_case_catcode:Nn</code>	<code>*</code>	<code>\token_case_meaning:NnTF</code>	$\langle test\ token \rangle$
<code>\token_case_catcode:NnTF</code>	<code>*</code>		{
<code>\token_case_charcode:Nn</code>	<code>*</code>		$\langle token\ case_1 \rangle$ $\{\langle code\ case_1 \rangle\}$
<code>\token_case_charcode:NnTF</code>	<code>*</code>		$\langle token\ case_2 \rangle$ $\{\langle code\ case_2 \rangle\}$
<code>\token_case_meaning:Nn</code>	<code>*</code>		...
<code>\token_case_meaning:NnTF</code>	<code>*</code>		$\langle token\ case_n \rangle$ $\{\langle code\ case_n \rangle\}$
			}
			$\{\langle true\ code \rangle\}$
			$\{\langle false\ code \rangle\}$

New: 2020-12-03

This function compares the $\langle test\ token \rangle$ in turn with each of the $\langle token\ cases \rangle$. If the two are equal (as described for `\token_if_eq_catcode:NNTF`, `\token_if_eq_charcode:NNTF` and `\token_if_eq_meaning:NNTF`, respectively) then the associated $\langle code \rangle$ is left in the input stream and other cases are discarded. If any of the cases are matched, the $\langle true\ code \rangle$ is also inserted into the input stream (after the code for the appropriate case), while if none match then the $\langle false\ code \rangle$ is inserted. The functions `\token_case_catcode:Nn`, `\token_case_charcode:Nn`, and `\token_case_meaning:Nn`, which do nothing if there is no match, are also available.

6 Peeking ahead at the next token

There is often a need to look ahead at the next token in the input stream while leaving it in place. This is handled using the “peek” functions. The generic `\peek_after:Nw` is provided along with a family of predefined tests for common cases. As peeking ahead does *not* skip spaces the predefined tests include both a space-respecting and space-skipping version. In addition, using `\peek_analysis_map_inline:n`, one can map through the following tokens in the input stream and repeatedly perform some tests.

<code>\peek_after:Nw</code>	<code>\peek_after:Nw</code>	$\langle function \rangle$	$\langle token \rangle$
-----------------------------	-----------------------------	----------------------------	-------------------------

Locally sets the test variable `\l_peek_token` equal to $\langle token \rangle$ (as an implicit token, *not* as a token list), and then expands the $\langle function \rangle$. The $\langle token \rangle$ remains in the input stream as the next item after the $\langle function \rangle$. The $\langle token \rangle$ here may be \sqcup , $\{$ or $\}$ (assuming normal T_EX category codes), *i.e.* it is not necessarily the next argument which would be grabbed by a normal function.

<u><u>\peek_gafter:Nw</u></u>	<u><u>\peek_gafter:Nw</u></u> $\langle function \rangle$ $\langle token \rangle$
	Globally sets the test variable <code>\g_peek_token</code> equal to $\langle token \rangle$ (as an implicit token, <i>not</i> as a token list), and then expands the $\langle function \rangle$. The $\langle token \rangle$ remains in the input stream as the next item after the $\langle function \rangle$. The $\langle token \rangle$ here may be \sqcup , $\{$ or $\}$ (assuming normal \TeX category codes), <i>i.e.</i> it is not necessarily the next argument which would be grabbed by a normal function.
<u><u>\l_peek_token</u></u>	Token set by <code>\peek_after:Nw</code> and available for testing as described above.
<u><u>\g_peek_token</u></u>	Token set by <code>\peek_gafter:Nw</code> and available for testing as described above.
<u><u>\peek_catcode:NTF</u></u>	<u><u>\peek_catcode:NTF</u></u> $\langle test\ token \rangle$ $\{\langle true\ code \rangle\}$ $\{\langle false\ code \rangle\}$
Updated: 2012-12-20	Tests if the next $\langle token \rangle$ in the input stream has the same category code as the $\langle test\ token \rangle$ (as defined by the test <code>\token_if_eq_catcode:NNTF</code>). Spaces are respected by the test and the $\langle token \rangle$ is left in the input stream after the $\langle true\ code \rangle$ or $\langle false\ code \rangle$ (as appropriate to the result of the test).
<u><u>\peek_catcode_ignore_spaces:NTF</u></u>	<u><u>\peek_catcode_ignore_spaces:NTF</u></u> $\langle test\ token \rangle$ $\{\langle true\ code \rangle\}$ $\{\langle false\ code \rangle\}$
Updated: 2012-12-20	Tests if the next non-space $\langle token \rangle$ in the input stream has the same category code as the $\langle test\ token \rangle$ (as defined by the test <code>\token_if_eq_catcode:NNTF</code>). Explicit and implicit space tokens (with character code 32 and category code 10) are ignored and removed by the test and the $\langle token \rangle$ is left in the input stream after the $\langle true\ code \rangle$ or $\langle false\ code \rangle$ (as appropriate to the result of the test).
<u><u>\peek_catcode_remove:NTF</u></u>	<u><u>\peek_catcode_remove:NTF</u></u> $\langle test\ token \rangle$ $\{\langle true\ code \rangle\}$ $\{\langle false\ code \rangle\}$
Updated: 2012-12-20	Tests if the next $\langle token \rangle$ in the input stream has the same category code as the $\langle test\ token \rangle$ (as defined by the test <code>\token_if_eq_catcode:NNTF</code>). Spaces are respected by the test and the $\langle token \rangle$ is removed from the input stream if the test is true. The function then places either the $\langle true\ code \rangle$ or $\langle false\ code \rangle$ in the input stream (as appropriate to the result of the test).
<u><u>\peek_catcode_remove_ignore_spaces:NTF</u></u>	<u><u>\peek_catcode_remove_ignore_spaces:NTF</u></u> $\langle test\ token \rangle$ $\{\langle true\ code \rangle\}$ $\{\langle false\ code \rangle\}$
Updated: 2012-12-20	Tests if the next non-space $\langle token \rangle$ in the input stream has the same category code as the $\langle test\ token \rangle$ (as defined by the test <code>\token_if_eq_catcode:NNTF</code>). Explicit and implicit space tokens (with character code 32 and category code 10) are ignored and removed by the test and the $\langle token \rangle$ is removed from the input stream if the test is true. The function then places either the $\langle true\ code \rangle$ or $\langle false\ code \rangle$ in the input stream (as appropriate to the result of the test).

<u>\peek_charcode:NTF</u>	\peek_charcode:NTF <test token> {(true code)} {(false code)}
Updated: 2012-12-20	Tests if the next <token> in the input stream has the same character code as the <test token> (as defined by the test \token_if_eq_charcode:NNTF). Spaces are respected by the test and the <token> is left in the input stream after the <true code> or <false code> (as appropriate to the result of the test).
<u>\peek_charcode_ignore_spaces:NTF</u>	\peek_charcode_ignore_spaces:NTF <test token> {(true code)} {(false code)}
Updated: 2012-12-20	Tests if the next non-space <token> in the input stream has the same character code as the <test token> (as defined by the test \token_if_eq_charcode:NNTF). Explicit and implicit space tokens (with character code 32 and category code 10) are ignored and removed by the test and the <token> is left in the input stream after the <true code> or <false code> (as appropriate to the result of the test).
<u>\peek_charcode_remove:NTF</u>	\peek_charcode_remove:NTF <test token> {(true code)} {(false code)}
Updated: 2012-12-20	Tests if the next <token> in the input stream has the same character code as the <test token> (as defined by the test \token_if_eq_charcode:NNTF). Spaces are respected by the test and the <token> is removed from the input stream if the test is true. The function then places either the <true code> or <false code> in the input stream (as appropriate to the result of the test).
<u>\peek_charcode_remove_ignore_spaces:NTF</u>	\peek_charcode_remove_ignore_spaces:NTF <test token> {(true code)} {(false code)}
Updated: 2012-12-20	Tests if the next non-space <token> in the input stream has the same character code as the <test token> (as defined by the test \token_if_eq_charcode:NNTF). Explicit and implicit space tokens (with character code 32 and category code 10) are ignored and removed by the test and the <token> is removed from the input stream if the test is true. The function then places either the <true code> or <false code> in the input stream (as appropriate to the result of the test).
<u>\peek_meaning:NTF</u>	\peek_meaning:NTF <test token> {(true code)} {(false code)}
Updated: 2011-07-02	Tests if the next <token> in the input stream has the same meaning as the <test token> (as defined by the test \token_if_eq_meaning:NNTF). Spaces are respected by the test and the <token> is left in the input stream after the <true code> or <false code> (as appropriate to the result of the test).
<u>\peek_meaning_ignore_spaces:NTF</u>	\peek_meaning_ignore_spaces:NTF <test token> {(true code)} {(false code)}
Updated: 2012-12-05	Tests if the next non-space <token> in the input stream has the same meaning as the <test token> (as defined by the test \token_if_eq_meaning:NNTF). Explicit and implicit space tokens (with character code 32 and category code 10) are ignored and removed by the test and the <token> is left in the input stream after the <true code> or <false code> (as appropriate to the result of the test).

<u>\peek_meaning_remove:NTF</u>	\peek_meaning_remove:NTF $\langle test\ token \rangle$ $\{\langle true\ code \rangle\}$ $\{\langle false\ code \rangle\}$
Updated: 2011-07-02	Tests if the next $\langle token \rangle$ in the input stream has the same meaning as the $\langle test\ token \rangle$ (as defined by the test \token_if_eq_meaning:NNTF). Spaces are respected by the test and the $\langle token \rangle$ is removed from the input stream if the test is true. The function then places either the $\langle true\ code \rangle$ or $\langle false\ code \rangle$ in the input stream (as appropriate to the result of the test).
<u>\peek_meaning_remove_ignore_spaces:NTF</u>	\peek_meaning_remove_ignore_spaces:NTF $\langle test\ token \rangle$ $\{\langle true\ code \rangle\}$ $\{\langle false\ code \rangle\}$
Updated: 2012-12-05	Tests if the next non-space $\langle token \rangle$ in the input stream has the same meaning as the $\langle test\ token \rangle$ (as defined by the test \token_if_eq_meaning:NNTF). Explicit and implicit space tokens (with character code 32 and category code 10) are ignored and removed by the test and the $\langle token \rangle$ is removed from the input stream if the test is true. The function then places either the $\langle true\ code \rangle$ or $\langle false\ code \rangle$ in the input stream (as appropriate to the result of the test).
<u>\peek_N_type:TF</u>	\peek_N_type:TF $\{\langle true\ code \rangle\}$ $\{\langle false\ code \rangle\}$
Updated: 2012-12-20	Tests if the next $\langle token \rangle$ in the input stream can be safely grabbed as an N-type argument. The test is $\langle false \rangle$ if the next $\langle token \rangle$ is either an explicit or implicit begin-group or end-group token (with any character code), or an explicit or implicit space character (with character code 32 and category code 10), or an outer token (never used in L ^A T _E X3) and $\langle true \rangle$ in all other cases. Note that a $\langle true \rangle$ result ensures that the next $\langle token \rangle$ is a valid N-type argument. However, if the next $\langle token \rangle$ is for instance \c_space_token, the test takes the $\langle false \rangle$ branch, even though the next $\langle token \rangle$ is in fact a valid N-type argument. The $\langle token \rangle$ is left in the input stream after the $\langle true\ code \rangle$ or $\langle false\ code \rangle$ (as appropriate to the result of the test).

<code>\peek_analysis_map_inline:n</code>	<code>\peek_analysis_map_inline:n {<i><inline function></i>}</code>
--	---

New: 2020-12-03

Repeatedly removes one *<token>* from the input stream and applies the *<inline function>* to it, until `\peek_analysis_map_break:` is called. The *<inline function>* receives three arguments for each *<token>* in the input stream:

- *<tokens>*, which both o-expand and x-expand to the *<token>*. The detailed form of *<tokens>* may change in later releases.
- *<char code>*, a decimal representation of the character code of the *<token>*, -1 if it is a control sequence.
- *<catcode>*, a capital hexadecimal digit which denotes the category code of the *<token>* (0: control sequence, 1: begin-group, 2: end-group, 3: math shift, 4: alignment tab, 6: parameter, 7: superscript, 8: subscript, A: space, B: letter, C: other, D: active). This can be converted to an integer by writing "*<catcode>*".

These arguments are the same as for `\tl_analysis_map_inline:nn` defined in `l3tl-analysis`. The *<char code>* and *<catcode>* do not take the meaning of a control sequence or active character into account: for instance, upon encountering the token `\c_group_begin_token` in the input stream, `\peek_analysis_map_inline:n` calls the *<inline function>* with #1 being `\exp_not:n { \c_group_begin_token }` (with the current implementation), #2 being -1 , and #3 being 0, as for any other control sequence. In contrast, upon encountering an explicit begin-group token `{`, the *<inline function>* is called with arguments `\exp_after:wN { \if_false: } \fi:`, 123 and 1.

The mapping is done at the current group level, *i.e.* any local assignments made by the *<inline function>* remain in effect after the loop. Within the code, `\l_peek_token` is set equal (as a token, not a token list) to the token under consideration.

<code>\peek_analysis_map_break:</code>	<code>\peek_analysis_map_inline:n</code>
<code>\peek_analysis_map_break:n</code>	<code>{ ... \peek_analysis_map_break:n {<i><code></i>} }</code>

New: 2020-12-03

Stops the `\peek_analysis_map_inline:n` loop from seeking more tokens, and inserts *<code>* in the input stream (empty for `\peek_analysis_map_break:`).

<code>\peek_regex:nTF</code>	<code>\peek_regex:nTF {<i><regex></i>} {<i><true code></i>} {<i><false code></i>}</code>
<code>\peek_regex:NTF</code>	

New: 2020-12-03

Tests if the *<tokens>* that follow in the input stream match the *<regular expression>*. Any *<tokens>* that have been read are left in the input stream after the *<true code>* or *<false code>* (as appropriate to the result of the test). See `l3regex` for documentation of the syntax of regular expressions. The *<regular expression>* is implicitly anchored at the start, so for instance `\peek_regex:nTF { a }` is essentially equivalent to `\peek_charcode:NTF a`.

T_EXhackers note: Implicit character tokens are correctly considered by `\peek_regex:nTF` as control sequences, while functions that inspect individual tokens (for instance `\peek_charcode:NTF`) only take into account their meaning.

```
\peek_regex_remove_once:nTF \peek_regex_remove_once:nTF {\<regex>} {\<true code>} {\<false code>}
\peek_regex_remove_once:NTF
```

New: 2020-12-03

Tests if the $\langle tokens \rangle$ that follow in the input stream match the $\langle regex \rangle$. If the test is true, the $\langle tokens \rangle$ are removed from the input stream and the $\langle true code \rangle$ is inserted, while if the test is false, the $\langle false code \rangle$ is inserted followed by the $\langle tokens \rangle$ that were originally in the input stream. See `l3regex` for documentation of the syntax of regular expressions. The $\langle regular expression \rangle$ is implicitly anchored at the start, so for instance `\peek_regex_remove_once:nTF { a }` is essentially equivalent to `\peek_charcode_remove:NTF a`.

TeXhackers note: Implicit character tokens are correctly considered by `\peek_regex_remove_once:nTF` as control sequences, while functions that inspect individual tokens (for instance `\peek_charcode:NTF`) only take into account their meaning.

```
\peek_regex_replace_once:nn \peek_regex_replace_once:nnTF {\<regex>} {\<replacement>} {\<true code>}
\peek_regex_replace_once:nnTF {\<false code>}
\peek_regex_replace_once:Nn
\peek_regex_replace_once:NnTF
```

New: 2020-12-03

If the $\langle tokens \rangle$ that follow in the input stream match the $\langle regex \rangle$, replaces them according to the $\langle replacement \rangle$ as for `\regex_replace_once:nnN`, and leaves the result in the input stream, after the $\langle true code \rangle$. Otherwise, leaves $\langle false code \rangle$ followed by the $\langle tokens \rangle$ that were originally in the input stream, with no modifications. See `l3regex` for documentation of the syntax of regular expressions and of the $\langle replacement \rangle$: for instance `\0` in the $\langle replacement \rangle$ is replaced by the tokens that were matched in the input stream. The $\langle regular expression \rangle$ is implicitly anchored at the start. In contrast to `\regex_replace_once:nnN`, no error arises if the $\langle replacement \rangle$ leads to an unbalanced token list: the tokens are inserted into the input stream without issue.

TeXhackers note: Implicit character tokens are correctly considered by `\peek_regex_replace_once:nnTF` as control sequences, while functions that inspect individual tokens (for instance `\peek_charcode:NTF`) only take into account their meaning.

7 Description of all possible tokens

Let us end by reviewing every case that a given token can fall into. This section is quite technical and some details are only meant for completeness. We distinguish the meaning of the token, which controls the expansion of the token and its effect on TeX's state, and its shape, which is used when comparing token lists such as for delimited arguments. Two tokens of the same shape must have the same meaning, but the converse does not hold.

A token has one of the following shapes.

- A control sequence, characterized by the sequence of characters that constitute its name: for instance, `\use:n` is a five-letter control sequence.

- An active character token, characterized by its character code (between 0 and 1114111 for LuaTeX and XeTeX and less for other engines) and category code 13.
- A character token, characterized by its character code and category code (one of 1, 2, 3, 4, 6, 7, 8, 10, 11 or 12 whose meaning is described below).

There are also a few internal tokens. The following list may be incomplete in some engines.

- Expanding `\the\font` results in a token that looks identical to the command that was used to select the current font (such as `\tenrm`) but it differs from it in shape.
- A “frozen” `\relax`, which differs from the primitive in shape (but has the same meaning), is inserted when the closing `\fi` of a conditional is encountered before the conditional is evaluated.
- Expanding `\noexpand <token>` (when the `<token>` is expandable) results in an internal token, displayed (temporarily) as `\notexpanded: <token>`, whose shape coincides with the `<token>` and whose meaning differs from `\relax`.
- An `\outer endtemplate:` can be encountered when peeking ahead at the next token; this expands to another internal token, `end of alignment template`.
- Tricky programming might access a frozen `\endwrite`.
- Some frozen tokens can only be accessed in interactive sessions: `\cr`, `\right`, `\endgroup`, `\fi`, `\inaccessible`.
- In LuaTeX, there is also the strange case of “bytes” `~~~~~1100xy` where x, y are any two lowercase hexadecimal digits, so that the hexadecimal number ranges from `\text{110000}=1114112$` to `~$1100ff = 1114367`. These are used to output individual bytes to files, rather than UTF-8. For the purposes of token comparisons they behave like non-expandable primitive control sequences (*not characters*) whose `\meaning` is the `\character_` followed by the given byte. If this byte is in the range 80–ff this gives an “invalid utf-8 sequence” error: applying `\token_to_str:N` or `\token_to_meaning:N` to these tokens is unsafe. Unfortunately, they don’t seem to be detectable safely by any means except perhaps Lua code.

The meaning of a (non-active) character token is fixed by its category code (and character code) and cannot be changed. We call these tokens *explicit* character tokens. Category codes that a character token can have are listed below by giving a sample output of the TeX primitive `\meaning`, together with their L^AT_EX3 names and most common example:

- 1 begin-group character (`group_begin`, often `{`),
- 2 end-group character (`group_end`, often `}`),
- 3 math shift character (`math_toggle`, often `$`),
- 4 alignment tab character (`alignment`, often `&`),
- 6 macro parameter character (`parameter`, often `#`),
- 7 superscript character (`math_superscript`, often `^`),

- 8 subscript character (`math_subscript`, often `_`),
- 10 blank space (`space`, often character code 32),
- 11 the letter (`letter`, such as `A`),
- 12 the character (`other`, such as `0`).

Category code 13 (`active`) is discussed below. Input characters can also have several other category codes which do not lead to character tokens for later processing: 0 (`escape`), 5 (`end_line`), 9 (`ignore`), 14 (`comment`), and 15 (`invalid`).

The meaning of a control sequence or active character can be identical to that of any character token listed above (with any character code), and we call such tokens *implicit* character tokens. The meaning is otherwise in the following list:

- a macro, used in $\text{\LaTeX}3$ for most functions and some variables (`tl`, `fp`, `seq`, ...),
- a primitive such as `\def` or `\topmark`, used in $\text{\LaTeX}3$ for some functions,
- a register such as `\count123`, used in $\text{\LaTeX}3$ for the implementation of some variables (`int`, `dim`, ...),
- a constant integer such as `\char"56` or `\mathchar"121`,
- a font selection command,
- undefined.

Macros can be `\protected` or not, `\long` or not (the opposite of what $\text{\LaTeX}3$ calls `nopar`), and `\outer` or not (unused in $\text{\LaTeX}3$). Their `\meaning` takes the form

$\langle prefix \rangle \text{ macro:} \langle argument \rangle \rightarrow \langle replacement \rangle$

where $\langle prefix \rangle$ is among `\protected`, `\long`, `\outer`, $\langle argument \rangle$ describes parameters that the macro expects, such as `#1#2#3`, and $\langle replacement \rangle$ describes how the parameters are manipulated, such as `\int_eval:n{\#2+\#1*\#3}`.

Now is perhaps a good time to mention some subtleties relating to tokens with category code 10 (space). Any input character with this category code (normally, space and tab characters) becomes a normal space, with character code 32 and category code 10.

When a macro takes an undelimited argument, explicit space characters (with character code 32 and category code 10) are ignored. If the following token is an explicit character token with category code 1 (begin-group) and an arbitrary character code, then \TeX scans ahead to obtain an equal number of explicit character tokens with category code 1 (begin-group) and 2 (end-group), and the resulting list of tokens (with outer braces removed) becomes the argument. Otherwise, a single token is taken as the argument for the macro: we call such single tokens “N-type”, as they are suitable to be used as an argument for a function with the signature `:N`.

When a macro takes a delimited argument \TeX scans ahead until finding the delimiter (outside any pairs of begin-group/end-group explicit characters), and the resulting list of tokens (with outer braces removed) becomes the argument. Note that explicit space characters at the start of the argument are *not* ignored in this case (and they prevent brace-stripping).

Part XVII

The l3prop package

Property lists

L^AT_EX3 implements a “property list” data type, which contain an unordered list of entries each of which consists of a $\langle key \rangle$ and an associated $\langle value \rangle$. The $\langle key \rangle$ and $\langle value \rangle$ may both be any $\langle balanced\ text \rangle$. It is possible to map functions to property lists such that the function is applied to every key–value pair within the list.

Each entry in a property list must have a unique $\langle key \rangle$: if an entry is added to a property list which already contains the $\langle key \rangle$ then the new entry overwrites the existing one. The $\langle keys \rangle$ are compared on a string basis, using the same method as `\str_if_eq:nn`.

Property lists are intended for storing key-based information for use within code. This is in contrast to key–value lists, which are a form of *input* parsed by the `keys` module.

1 Creating and initialising property lists

```
\prop_new:N
\prop_new:c
```

```
\prop_new:N <property list>
```

Creates a new $\langle property\ list \rangle$ or raises an error if the name is already taken. The declaration is global. The $\langle property\ list \rangle$ initially contains no entries.

```
\prop_clear:N
\prop_clear:c
\prop_gclear:N
\prop_gclear:c
```

```
\prop_clear:N <property list>
```

Clears all entries from the $\langle property\ list \rangle$.

```
\prop_clear_new:N
\prop_clear_new:c
\prop_gclear_new:N
\prop_gclear_new:c
```

```
\prop_clear_new:N <property list>
```

Ensures that the $\langle property\ list \rangle$ exists globally by applying `\prop_new:N` if necessary, then applies `\prop_(g)clear:N` to leave the list empty.

```
\prop_set_eq:NN
\prop_set_eq:(cN|Nc|cc)
\prop_gset_eq:NN
\prop_gset_eq:(cN|Nc|cc)
```

```
\prop_set_eq:NN <property list1> <property list2>
```

Sets the content of $\langle property\ list_1 \rangle$ equal to that of $\langle property\ list_2 \rangle$.

```
\prop_set_from_keyval:Nn
\prop_set_from_keyval:cn
\prop_gset_from_keyval:Nn
\prop_gset_from_keyval:cn
```

```
\prop_set_from_keyval:Nn <prop var>
```

```
{
  <key1> = <value1> ,
  <key2> = <value2> , ...
}
```

Sets $\langle prop\ var \rangle$ to contain key–value pairs given in the second argument. If duplicate keys appear only one of the values is kept.

New: 2017-11-28
Updated: 2019-08-25

```
\prop_const_from_keyval:Nn
\prop_const_from_keyval:cn
```

New: 2017-11-28
Updated: 2019-08-25

```
\prop_const_from_keyval:Nn <prop var>
{
  <key1> = <value1> ,
  <key2> = <value2> , ...
}
```

Creates a new constant *<prop var>* or raises an error if the name is already taken. The *<prop var>* is set globally to contain key–value pairs given in the second argument. If duplicate keys appear only one of the values is kept.

2 Adding entries to property lists

```
\prop_put:Nnn
\prop_put:(NnV|Nno|Nnx|NVn|NVV|NVx|Nvx|Non|Noo|Nxx|cnn|cnV|cno|cnx|cVn|cVV|cVx|cvx|con|coo|cxx)
\prop_gput:Nnn
\prop_gput:(NnV|Nno|Nnx|NVn|NVV|NVx|Nvx|Non|Noo|Nxx|cnn|cnV|cno|cnx|cVn|cVV|cVx|cvx|con|coo|cxx)
```

```
\prop_put:Nnn
  <property list>
  {<key>}
  {<value>}
```

Updated: 2012-07-09

Adds an entry to the *<property list>* which may be accessed using the *<key>* and which has *<value>*. Both the *<key>* and *<value>* may contain any *<balanced text>*. The *<key>* is stored after processing with `\tl_to_str:n`, meaning that category codes are ignored. If the *<key>* is already present in the *<property list>*, the existing entry is overwritten by the new *<value>*.

```
\prop_put_if_new:Nnn
\prop_put_if_new:cnn
\prop_gput_if_new:Nnn
\prop_gput_if_new:cnn
```

```
\prop_put_if_new:Nnn <property list> {<key>} {<value>}
```

If the *<key>* is present in the *<property list>* then no action is taken. If the *<key>* is not present in the *<property list>* then a new entry is added. Both the *<key>* and *<value>* may contain any *<balanced text>*. The *<key>* is stored after processing with `\tl_to_str:n`, meaning that category codes are ignored.

3 Recovering values from property lists

```
\prop_get:NnN
\prop_get:(NVN|NvN|NoN|cnN|cVN|cvN|coN)
```

```
\prop_get:NnN <property list> {<key>} <tl var>
```

Updated: 2011-08-28

Recovers the *<value>* stored with *<key>* from the *<property list>*, and places this in the *<token list variable>*. If the *<key>* is not found in the *<property list>* then the *<token list variable>* is set to the special marker `\q_no_value`. The *<token list variable>* is set within the current T_EX group. See also `\prop_get:NnNTF`.

```
\prop_pop:NnN
\prop_pop:(NoN|cnN|coN)
```

Updated: 2011-08-18

```
\prop_pop:NnN <property list> {<key>} <tl var>
```

Recovers the *<value>* stored with *<key>* from the *<property list>*, and places this in the *<token list variable>*. If the *<key>* is not found in the *<property list>* then the *<token list variable>* is set to the special marker `\q_no_value`. The *<key>* and *<value>* are then deleted from the property list. Both assignments are local. See also `\prop_pop:NnNTF`.

<code>\prop_gpop:NnN</code>
<code>\prop_gpop:(NoN cnN coN)</code> ★
Updated: 2011-08-18

`\prop_gpop:NnN` $\langle property\ list \rangle$ $\{\langle key \rangle\}$ $\langle tl\ var \rangle$

Recovers the $\langle value \rangle$ stored with $\langle key \rangle$ from the $\langle property\ list \rangle$, and places this in the $\langle token\ list\ variable \rangle$. If the $\langle key \rangle$ is not found in the $\langle property\ list \rangle$ then the $\langle token\ list\ variable \rangle$ is set to the special marker `\q_no_value`. The $\langle key \rangle$ and $\langle value \rangle$ are then deleted from the property list. The $\langle property\ list \rangle$ is modified globally, while the assignment of the $\langle token\ list\ variable \rangle$ is local. See also `\prop_gpop:NnNTF`.

<code>\prop_item:Nn</code> ★
<code>\prop_item:cn</code> ★
New: 2014-07-17

`\prop_item:Nn` $\langle property\ list \rangle$ $\{\langle key \rangle\}$

Expands to the $\langle value \rangle$ corresponding to the $\langle key \rangle$ in the $\langle property\ list \rangle$. If the $\langle key \rangle$ is missing, this has an empty expansion.

TeXhackers note: This function is slower than the non-expandable analogue `\prop_get:NnN`. The result is returned within the `\unexpanded` primitive (`\exp_not:n`), which means that the $\langle value \rangle$ does not expand further when appearing in an x-type argument expansion.

<code>\prop_count:N</code> ★
<code>\prop_count:c</code> ★

`\prop_count:N` $\langle property\ list \rangle$

Leaves the number of key–value pairs in the $\langle property\ list \rangle$ in the input stream as an $\langle integer\ denotation \rangle$.

4 Modifying property lists

<code>\prop_remove:Nn</code>
<code>\prop_remove:(NV cn cV)</code>
<code>\prop_gremove:Nn</code>
<code>\prop_gremove:(NV cn cV)</code>
New: 2012-05-12

`\prop_remove:Nn` $\langle property\ list \rangle$ $\{\langle key \rangle\}$

Removes the entry listed under $\langle key \rangle$ from the $\langle property\ list \rangle$. If the $\langle key \rangle$ is not found in the $\langle property\ list \rangle$ no change occurs, *i.e.* there is no need to test for the existence of a key before deleting it.

5 Property list conditionals

<code>\prop_if_exist_p:N</code> ★
<code>\prop_if_exist_p:c</code> ★
<code>\prop_if_exist:NTF</code> ★
<code>\prop_if_exist:cTF</code> ★
New: 2012-03-03

`\prop_if_exist_p:N` $\langle property\ list \rangle$

`\prop_if_exist:NTF` $\langle property\ list \rangle$ $\{\langle true\ code \rangle\}$ $\{\langle false\ code \rangle\}$

Tests whether the $\langle property\ list \rangle$ is currently defined. This does not check that the $\langle property\ list \rangle$ really is a property list variable.

<code>\prop_if_empty_p:N</code> ★
<code>\prop_if_empty_p:c</code> ★
<code>\prop_if_empty:NTF</code> ★
<code>\prop_if_empty:cTF</code> ★

`\prop_if_empty_p:N` $\langle property\ list \rangle$

`\prop_if_empty:NTF` $\langle property\ list \rangle$ $\{\langle true\ code \rangle\}$ $\{\langle false\ code \rangle\}$

Tests if the $\langle property\ list \rangle$ is empty (containing no entries).

<code>\prop_if_in_p:Nn</code>	★	<code>\prop_if_in:NnTF</code>	$\langle property\ list \rangle$	$\{\langle key \rangle\}$	$\{\langle true\ code \rangle\}$	$\{\langle false\ code \rangle\}$
<code>\prop_if_in_p:(NV No cn cV co)</code>	★					
<code>\prop_if_in:NnTF</code>	★					
<code>\prop_if_in:(NV No cn cV co)TF</code>	★					

Updated: 2011-09-15

Tests if the $\langle key \rangle$ is present in the $\langle property\ list \rangle$, making the comparison using the method described by `\str_if_eq:NnTF`.

TeXhackers note: This function iterates through every key-value pair in the $\langle property\ list \rangle$ and is therefore slower than using the non-expandable `\prop_get:NnTF`.

6 Recovering values from property lists with branching

The functions in this section combine tests for the presence of a key in a property list with recovery of the associated valued. This makes them useful for cases where different cases follow dependent on the presence or absence of a key in a property list. They offer increased readability and performance over separate testing and recovery phases.

<code>\prop_get:NnTF</code>	<code>\prop_get:NnTF</code>	$\langle property\ list \rangle$	$\{\langle key \rangle\}$	$\langle token\ list\ variable \rangle$
<code>\prop_get:(NVN NvN NoN cnN cVN cvN coN)TF</code>			$\{\langle true\ code \rangle\}$	$\{\langle false\ code \rangle\}$

Updated: 2012-05-19

If the $\langle key \rangle$ is not present in the $\langle property\ list \rangle$, leaves the $\langle false\ code \rangle$ in the input stream. The value of the $\langle token\ list\ variable \rangle$ is not defined in this case and should not be relied upon. If the $\langle key \rangle$ is present in the $\langle property\ list \rangle$, stores the corresponding $\langle value \rangle$ in the $\langle token\ list\ variable \rangle$ without removing it from the $\langle property\ list \rangle$, then leaves the $\langle true\ code \rangle$ in the input stream. The $\langle token\ list\ variable \rangle$ is assigned locally.

<code>\prop_pop:NnTF</code>	<code>\prop_pop:NnTF</code>	$\langle property\ list \rangle$	$\{\langle key \rangle\}$	$\langle token\ list\ variable \rangle$	$\{\langle true\ code \rangle\}$
<code>\prop_pop:cnTF</code>					$\{\langle false\ code \rangle\}$

New: 2011-08-18
Updated: 2012-05-19

If the $\langle key \rangle$ is not present in the $\langle property\ list \rangle$, leaves the $\langle false\ code \rangle$ in the input stream. The value of the $\langle token\ list\ variable \rangle$ is not defined in this case and should not be relied upon. If the $\langle key \rangle$ is present in the $\langle property\ list \rangle$, pops the corresponding $\langle value \rangle$ in the $\langle token\ list\ variable \rangle$, *i.e.* removes the item from the $\langle property\ list \rangle$. Both the $\langle property\ list \rangle$ and the $\langle token\ list\ variable \rangle$ are assigned locally.

<code>\prop_gpop:NnTF</code>	<code>\prop_gpop:NnTF</code>	$\langle property\ list \rangle$	$\{\langle key \rangle\}$	$\langle token\ list\ variable \rangle$	$\{\langle true\ code \rangle\}$
<code>\prop_gpop:cnTF</code>					$\{\langle false\ code \rangle\}$

New: 2011-08-18
Updated: 2012-05-19

If the $\langle key \rangle$ is not present in the $\langle property\ list \rangle$, leaves the $\langle false\ code \rangle$ in the input stream. The value of the $\langle token\ list\ variable \rangle$ is not defined in this case and should not be relied upon. If the $\langle key \rangle$ is present in the $\langle property\ list \rangle$, pops the corresponding $\langle value \rangle$ in the $\langle token\ list\ variable \rangle$, *i.e.* removes the item from the $\langle property\ list \rangle$. The $\langle property\ list \rangle$ is modified globally, while the $\langle token\ list\ variable \rangle$ is assigned locally.

7 Mapping to property lists

All mappings are done at the current group level, *i.e.* any local assignments made by the $\langle function \rangle$ or $\langle code \rangle$ discussed below remain in effect after the loop.

$\backslash prop_map_function:Nn$	☆
$\backslash prop_map_function:cN$	☆
Updated: 2013-01-08	

$\backslash prop_map_function:Nn$ $\langle property\ list \rangle$ $\langle function \rangle$

Applies $\langle function \rangle$ to every $\langle entry \rangle$ stored in the $\langle property\ list \rangle$. The $\langle function \rangle$ receives two arguments for each iteration: the $\langle key \rangle$ and associated $\langle value \rangle$. The order in which $\langle entries \rangle$ are returned is not defined and should not be relied upon. To pass further arguments to the $\langle function \rangle$, see $\backslash prop_map_tokens:Nn$.

$\backslash prop_map_inline:Nn$	
$\backslash prop_map_inline:cN$	
Updated: 2013-01-08	

$\backslash prop_map_inline:Nn$ $\langle property\ list \rangle$ $\{ \langle inline\ function \rangle \}$

Applies $\langle inline\ function \rangle$ to every $\langle entry \rangle$ stored within the $\langle property\ list \rangle$. The $\langle inline\ function \rangle$ should consist of code which receives the $\langle key \rangle$ as #1 and the $\langle value \rangle$ as #2. The order in which $\langle entries \rangle$ are returned is not defined and should not be relied upon.

$\backslash prop_map_tokens:Nn$	☆
$\backslash prop_map_tokens:cN$	☆

$\backslash prop_map_tokens:Nn$ $\langle property\ list \rangle$ $\{ \langle code \rangle \}$

Analogue of $\backslash prop_map_function:Nn$ which maps several tokens instead of a single function. The $\langle code \rangle$ receives each key–value pair in the $\langle property\ list \rangle$ as two trailing brace groups. For instance,

```
 $\backslash prop\_map\_tokens:Nn \backslash l\_my\_prop \{ \backslash str\_if\_eq:nnT \{ mykey \} \}$ 
```

expands to the value corresponding to `mykey`: for each pair in $\backslash l_my_prop$ the function $\backslash str_if_eq:nnT$ receives `mykey`, the $\langle key \rangle$ and the $\langle value \rangle$ as its three arguments. For that specific task, $\backslash prop_item:Nn$ is faster.

$\backslash prop_map_break:$	☆
Updated: 2012-06-29	

$\backslash prop_map_break:$

Used to terminate a $\backslash prop_map_...$ function before all entries in the $\langle property\ list \rangle$ have been processed. This normally takes place within a conditional statement, for example

```
 $\backslash prop\_map\_inline:Nn \backslash l\_my\_prop$ 
{
   $\backslash str\_if\_eq:nnTF \{ \#1 \} \{ bingo \}$ 
  {  $\backslash prop\_map\_break:$  }
  {
    % Do something useful
  }
}
```

Use outside of a $\backslash prop_map_...$ scenario leads to low level T_EX errors.

T_EXhackers note: When the mapping is broken, additional tokens may be inserted before further items are taken from the input stream. This depends on the design of the mapping function.

<hr/> <code>\prop_map_break:n</code> ☆ <hr/>	<code>\prop_map_break:n {<code>}</code>
Updated: 2012-06-29 <hr/>	Used to terminate a <code>\prop_map_...</code> function before all entries in the <i><property list></i> have been processed, inserting the <i><code></i> after the mapping has ended. This normally takes place within a conditional statement, for example

```

\prop_map_inline:Nn \l_my_prop
{
  \str_if_eq:nnTF { #1 } { bingo }
  { \prop_map_break:n { <code> } }
  {
    % Do something useful
  }
}

```

Use outside of a `\prop_map_...` scenario leads to low level T_EX errors.

T_EXhackers note: When the mapping is broken, additional tokens may be inserted before the *<code>* is inserted into the input stream. This depends on the design of the mapping function.

8 Viewing property lists

<hr/> <code>\prop_show:N</code> <code>\prop_show:c</code> <hr/>	<code>\prop_show:N <property list></code>
Updated: 2015-08-01 <hr/>	Displays the entries in the <i><property list></i> in the terminal.

<hr/> <code>\prop_log:N</code> <code>\prop_log:c</code> <hr/>	<code>\prop_log:N <property list></code>
New: 2014-08-12 Updated: 2015-08-01 <hr/>	Writes the entries in the <i><property list></i> in the log file.

9 Scratch property lists

<hr/> <code>\l_tmpa_prop</code> <code>\l_tmpb_prop</code> <hr/>	Scratch property lists for local assignment. These are never used by the kernel code, and so are safe for use with any L ^A T _E X3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.
New: 2012-06-23 <hr/>	

<hr/> <code>\g_tmpa_prop</code> <code>\g_tmpb_prop</code> <hr/>	Scratch property lists for global assignment. These are never used by the kernel code, and so are safe for use with any L ^A T _E X3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.
New: 2012-06-23 <hr/>	

10 Constants

<u><u>\c_empty_prop</u></u>	A permanently-empty property list used for internal comparisons.
-----------------------------	--

Part XVIII

The l3msg package

Messages

Messages need to be passed to the user by modules, either when errors occur or to indicate how the code is proceeding. The `l3msg` module provides a consistent method for doing this (as opposed to writing directly to the terminal or log).

The system used by `l3msg` to create messages divides the process into two distinct parts. Named messages are created in the first part of the process; at this stage, no decision is made about the type of output that the message will produce. The second part of the process is actually producing a message. At this stage a choice of message *class* has to be made, for example `error`, `warning` or `info`.

By separating out the creation and use of messages, several benefits are available. First, the messages can be altered later without needing details of where they are used in the code. This makes it possible to alter the language used, the detail level and so on. Secondly, the output which results from a given message can be altered. This can be done on a message class, module or message name basis. In this way, message behaviour can be altered and messages can be entirely suppressed.

1 Creating new messages

All messages have to be created before they can be used. The text of messages is automatically wrapped to the length available in the console. As a result, formatting is only needed where it helps to show meaning. In particular, `\\` may be used to force a new line and `_` forces an explicit space. Additionally, `\{`, `\#`, `\}`, `\%` and `\~` can be used to produce the corresponding character.

Messages may be subdivided *by one level* using the `/` character. This is used within the message filtering system to allow for example the L^AT_EX kernel messages to belong to the module `LaTeX` while still being filterable at a more granular level. Thus for example

```
\msg_new:nnnn { mymodule } { submodule / message } ...
```

will allow to filter out specifically messages from the `submodule`.

```
\msg_new:nnnn
\msg_new:nnn
Updated: 2011-08-16
```

```
\msg_new:nnnn {<module>} {<message>} {<text>} {<more text>}
```

Creates a `<message>` for a given `<module>`. The message is defined to first give `<text>` and then `<more text>` if the user requests it. If no `<more text>` is available then a standard text is given instead. Within `<text>` and `<more text>` four parameters (`#1` to `#4`) can be used: these will be supplied at the time the message is used. An error is raised if the `<message>` already exists.

```
\msg_set:nnnn
\msg_set:nnn
\msg_gset:nnnn
\msg_gset:nnn
```

```
\msg_set:nnnn {<module>} {<message>} {<text>} {<more text>}
```

Sets up the text for a `<message>` for a given `<module>`. The message is defined to first give `<text>` and then `<more text>` if the user requests it. If no `<more text>` is available then a standard text is given instead. Within `<text>` and `<more text>` four parameters (`#1` to `#4`) can be used: these will be supplied at the time the message is used.

<hr/>	<hr/>
<code>\msg_if_exist_p:nn *</code>	<code>\msg_if_exist_p:nn {<module>} {<message>}</code>
<code>\msg_if_exist:nnTF *</code>	<code>\msg_if_exist:nnTF {<module>} {<message>} {<true code>} {<false code>}</code>
<hr/>	<hr/>
New: 2012-03-03	Tests whether the <i><message></i> for the <i><module></i> is currently defined.

2 Contextual information for messages

<hr/>	<hr/>
<code>\msg_line_context: ☆</code>	<code>\msg_line_context:</code>
<hr/>	<hr/>
	Prints the current line number when a message is given, and thus suitable for giving context to messages. The number itself is proceeded by the text on line.

<hr/>	<hr/>
<code>\msg_line_number: *</code>	<code>\msg_line_number:</code>
<hr/>	<hr/>
	Prints the current line number when a message is given.

<hr/>	<hr/>
<code>\msg_fatal_text:n *</code>	<code>\msg_fatal_text:n {<module>}</code>
<hr/>	<hr/>
	Produces the standard text
	Fatal Package <i><module></i> Error
	This function can be redefined to alter the language in which the message is given, using #1 as the name of the <i><module></i> to be included.

<hr/>	<hr/>
<code>\msg_critical_text:n *</code>	<code>\msg_critical_text:n {<module>}</code>
<hr/>	<hr/>
	Produces the standard text
	Critical Package <i><module></i> Error
	This function can be redefined to alter the language in which the message is given, using #1 as the name of the <i><module></i> to be included.

<hr/>	<hr/>
<code>\msg_error_text:n *</code>	<code>\msg_error_text:n {<module>}</code>
<hr/>	<hr/>
	Produces the standard text
	Package <i><module></i> Error
	This function can be redefined to alter the language in which the message is given, using #1 as the name of the <i><module></i> to be included.

<hr/>	<hr/>
<code>\msg_warning_text:n *</code>	<code>\msg_warning_text:n {<module>}</code>
<hr/>	<hr/>
	Produces the standard text
	Package <i><module></i> Warning
	This function can be redefined to alter the language in which the message is given, using #1 as the name of the <i><module></i> to be included. The <i><type></i> of <i><module></i> may be adjusted: Package is the standard outcome: see <code>\msg_module_type:n</code> .

<hr/> <code>\msg_info_text:n</code> ★ <hr/>	<code>\msg_info_text:n {⟨module⟩}</code> Produces the standard text: <div>Package <code>⟨module⟩</code> Info</div> This function can be redefined to alter the language in which the message is given, using #1 as the name of the <code>⟨module⟩</code> to be included. The <code>⟨type⟩</code> of <code>⟨module⟩</code> may be adjusted: <code>Package</code> is the standard outcome: see <code>\msg_module_type:n</code> .
<hr/> <code>\msg_module_name:n</code> ★ <hr/> <div>New: 2018-10-10</div>	<code>\msg_module_name:n {⟨module⟩}</code> Expands to the public name of the <code>⟨module⟩</code> as defined by <code>\g_msg_module_name_prop</code> (or otherwise leaves the <code>⟨module⟩</code> unchanged).
<hr/> <code>\msg_module_type:n</code> ★ <hr/> <div>New: 2018-10-10</div>	<code>\msg_module_type:n {⟨module⟩}</code> Expands to the description which applies to the <code>⟨module⟩</code> , for example a <code>Package</code> or <code>Class</code> . The information here is defined in <code>\g_msg_module_type_prop</code> , and will default to <code>Package</code> if an entry is not present.
<hr/> <code>\msg_see_documentation_text:n</code> ★ <hr/> <div>Updated: 2018-09-30</div>	<code>\msg_see_documentation_text:n {⟨module⟩}</code> Produces the standard text <div>See the <code>⟨module⟩</code> documentation for further information.</div> This function can be redefined to alter the language in which the message is given, using #1 as the name of the <code>⟨module⟩</code> to be included. The name of the <code>⟨module⟩</code> may be altered by use of <code>\g_msg_module_documentation_prop</code>
<hr/> <code>\g_msg_module_name_prop</code> <hr/> <div>New: 2018-10-10</div>	Provides a mapping between the module name used for messages, and that for documentation. For example, <code>L^AT_EX3</code> core messages are stored in the reserved <code>L^AT_EX</code> tree, but are printed as <code>L^AT_EX3</code> .
<hr/> <code>\g_msg_module_type_prop</code> <hr/> <div>New: 2018-10-10</div>	Provides a mapping between the module name used for messages, and that type of module. For example, for <code>L^AT_EX3</code> core messages, an empty entry is set here meaning that they are not described using the standard <code>Package</code> text.

3 Issuing messages

Messages behave differently depending on the message class. In all cases, the message may be issued supplying 0 to 4 arguments. If the number of arguments supplied here does not match the number in the definition of the message, extra arguments are ignored, or empty arguments added (of course the sense of the message may be impaired). The four arguments are converted to strings before being added to the message text: the `x`-type variants should be used to expand material.

```

\msg_fatal:nnnnnn
\msg_fatal:nnxxxx
\msg_fatal:nnnnn
\msg_fatal:nnxxx
\msg_fatal:nnnn
\msg_fatal:nnxx
\msg_fatal:nnn
\msg_fatal:nnx
\msg_fatal:nn

```

Updated: 2012-08-11

```

\msg_fatal:nnnnnn {\module} {\message} {\arg one} {\arg two} {\arg three}
{\arg four}

```

Issues $\langle module \rangle$ error $\langle message \rangle$, passing $\langle arg one \rangle$ to $\langle arg four \rangle$ to the text-creating functions. After issuing a fatal error the T_EX run halts. No PDF file will be produced in this case (DVI mode runs may produce a truncated DVI file).

```

\msg_critical:nnnnnn
\msg_critical:nnxxxx
\msg_critical:nnnnn
\msg_critical:nnxxx
\msg_critical:nnnn
\msg_critical:nnxx
\msg_critical:nnn
\msg_critical:nnx
\msg_critical:nn

```

Updated: 2012-08-11

```

\msg_critical:nnnnnn {\module} {\message} {\arg one} {\arg two} {\arg three}
{\arg four}

```

Issues $\langle module \rangle$ error $\langle message \rangle$, passing $\langle arg one \rangle$ to $\langle arg four \rangle$ to the text-creating functions. After issuing a critical error, T_EX stops reading the current input file. This may halt the T_EX run (if the current file is the main file) or may abort reading a sub-file.

T_EXhackers note: The T_EX `\endinput` primitive is used to exit the file. In particular, the rest of the current line remains in the input stream.

```

\msg_error:nnnnnn
\msg_error:nnxxxx
\msg_error:nnnnn
\msg_error:nnxxx
\msg_error:nnnn
\msg_error:nnxx
\msg_error:nnn
\msg_error:nnx
\msg_error:nn

```

Updated: 2012-08-11

```

\msg_error:nnnnnn {\module} {\message} {\arg one} {\arg two} {\arg three}
{\arg four}

```

Issues $\langle module \rangle$ error $\langle message \rangle$, passing $\langle arg one \rangle$ to $\langle arg four \rangle$ to the text-creating functions. The error interrupts processing and issues the text at the terminal. After user input, the run continues.

```

\msg_warning:nnnnnn
\msg_warning:nnxxxx
\msg_warning:nnnnn
\msg_warning:nnxxx
\msg_warning:nnnn
\msg_warning:nnxx
\msg_warning:nnn
\msg_warning:nnx
\msg_warning:nn

```

Updated: 2012-08-11

```

\msg_warning:nnxxxx {\module} {\message} {\arg one} {\arg two} {\arg three}
{\arg four}

```

Issues $\langle module \rangle$ warning $\langle message \rangle$, passing $\langle arg one \rangle$ to $\langle arg four \rangle$ to the text-creating functions. The warning text is added to the log file and the terminal, but the T_EX run is not interrupted.

<hr/> <pre> \msg_info:nnnnnn \msg_info:nnxxxx \msg_info:nnnnn \msg_info:nnxxx \msg_info:nnnn \msg_info:nnxx \msg_info:nnn \msg_info:nnx \msg_info:nn </pre> <hr/> <p>Updated: 2012-08-11</p> <hr/>	<pre> \msg_info:nnnnnn {\module} {\message} {\arg one} {\arg two} {\arg three} {\arg four} </pre> <p>Issues <i>⟨module⟩</i> information <i>⟨message⟩</i>, passing <i>⟨arg one⟩</i> to <i>⟨arg four⟩</i> to the text-creating functions. The information text is added to the log file.</p>
<hr/> <pre> \msg_log:nnnnnn \msg_log:nnxxxx \msg_log:nnnnn \msg_log:nnxxx \msg_log:nnnn \msg_log:nnxx \msg_log:nnn \msg_log:nnx \msg_log:nn </pre> <hr/> <p>Updated: 2012-08-11</p> <hr/>	<pre> \msg_log:nnnnnn {\module} {\message} {\arg one} {\arg two} {\arg three} {\arg four} </pre> <p>Issues <i>⟨module⟩</i> information <i>⟨message⟩</i>, passing <i>⟨arg one⟩</i> to <i>⟨arg four⟩</i> to the text-creating functions. The information text is added to the log file: the output is briefer than <code>\msg_info:nnnnnn</code>.</p>
<hr/> <pre> \msg_term:nnnnnn \msg_term:nnxxxx \msg_term:nnnnn \msg_term:nnxxx \msg_term:nnnn \msg_term:nnxx \msg_term:nnn \msg_term:nnx \msg_term:nn </pre> <hr/> <p>New: 2020-07-16</p> <hr/>	<pre> \msg_term:nnnnnn {\module} {\message} {\arg one} {\arg two} {\arg three} {\arg four} </pre> <p>Issues <i>⟨module⟩</i> information <i>⟨message⟩</i>, passing <i>⟨arg one⟩</i> to <i>⟨arg four⟩</i> to the text-creating functions. The information text is printed on the terminal (and added to the log file): the output is similar to that of <code>\msg_log:nnnnnn</code>.</p>
<hr/> <pre> \msg_none:nnnnnn \msg_none:nnxxxx \msg_none:nnnnn \msg_none:nnxxx \msg_none:nnnn \msg_none:nnxx \msg_none:nnn \msg_none:nnx \msg_none:nn </pre> <hr/> <p>Updated: 2012-08-11</p> <hr/>	<pre> \msg_none:nnnnnn {\module} {\message} {\arg one} {\arg two} {\arg three} {\arg four} </pre> <p>Does nothing: used as a message class to prevent any output at all (see the discussion of message redirection).</p>

3.1 Expandable error messages

In very rare cases it may be necessary to produce errors in an expansion-only context. The functions in this section should only be used if there is no alternative approach using `\msg_error:nnnnnn` or other non-expandable commands from the previous section.

Despite having a similar interface as non-expandable messages, expandable errors must be handled internally very differently from normal error messages, as none of the tools to print to the terminal or the log file are expandable. As a result, short-hands such as `\{` or `\|` do not work, and messages must be very short (with default settings, they are truncated after approximately 50 characters). It is advisable to ensure that the message is understandable even when truncated, by putting the most important information up front. Another particularity of expandable messages is that they cannot be redirected or turned off by the user.

<code>\msg_expandable_error:nnnnnn</code>	*	<code>\msg_expandable_error:nnnnnn {<module>} {<message>} {<arg one>} {<arg</code>
<code>\msg_expandable_error:nnffff</code>	*	<code>two)} {<arg three>} {<arg four>}</code>
<code>\msg_expandable_error:nnnnn</code>	*	
<code>\msg_expandable_error:nnfff</code>	*	
<code>\msg_expandable_error:nnnn</code>	*	
<code>\msg_expandable_error:nnff</code>	*	
<code>\msg_expandable_error:nnn</code>	*	
<code>\msg_expandable_error:nnf</code>	*	
<code>\msg_expandable_error:nn</code>	*	

New: 2015-08-06

Updated: 2019-02-28

Issues an “Undefined error” message from T_EX itself using the undefined control sequence `\::error` then prints “! *<module>*: ”*<error message>*”, which should be short. With default settings, anything beyond approximately 60 characters long (or bytes in some engines) is cropped. A leading space might be removed as well.

4 Redirecting messages

Each message has a “name”, which can be used to alter the behaviour of the message when it is given. Thus we might have

```
\msg_new:nnnn { module } { my-message } { Some-text } { Some-more-text }
```

to define a message, with

```
\msg_error:nn { module } { my-message }
```

when it is used. With no filtering, this raises an error. However, we could alter the behaviour with

```
\msg_redirect_class:nn { error } { warning }
```

to turn all errors into warnings, or with

```
\msg_redirect_module:nnn { module } { error } { warning }
```

to alter only messages from that module, or even

```
\msg_redirect_name:nnn { module } { my-message } { warning }
```

to target just one message. Redirection applies first to individual messages, then to messages from one module and finally to messages of one class. Thus it is possible to select out an individual message for special treatment even if the entire class is already redirected.

Multiple redirections are possible. Redirections can be cancelled by providing an empty argument for the target class. Redirection to a missing class raises an error immediately. Infinite loops are prevented by eliminating the redirection starting from the target of the redirection that caused the loop to appear. Namely, if redirections are requested as $A \rightarrow B$, $B \rightarrow C$ and $C \rightarrow A$ in this order, then the $A \rightarrow B$ redirection is cancelled.

`\msg_redirect_class:nn`

Updated: 2012-04-27

`\msg_redirect_class:nn {<class one>} {<class two>}`

Changes the behaviour of messages of *<class one>* so that they are processed using the code for those of *<class two>*.

`\msg_redirect_module:nnn`

Updated: 2012-04-27

`\msg_redirect_module:nnn {<module>} {<class one>} {<class two>}`

Redirects message of *<class one>* for *<module>* to act as though they were from *<class two>*. Messages of *<class one>* from sources other than *<module>* are not affected by this redirection. This function can be used to make some messages “silent” by default. For example, all of the **warning** messages of *<module>* could be turned off with:

`\msg_redirect_module:nnn { module } { warning } { none }`

`\msg_redirect_name:nnn`

Updated: 2012-04-27

`\msg_redirect_name:nnn {<module>} {<message>} {<class>}`

Redirects a specific *<message>* from a specific *<module>* to act as a member of *<class>* of messages. No further redirection is performed. This function can be used to make a selected message “silent” without changing global parameters:

`\msg_redirect_name:nnn { module } { annoying-message } { none }`

Part XIX

The l3file package

File and I/O operations

This module provides functions for working with external files. Some of these functions apply to an entire file, and have prefix `\file_...`, while others are used to work with files on a line by line basis and have prefix `\ior_...` (reading) or `\iow_...` (writing).

It is important to remember that when reading external files \TeX attempts to locate them using both the operating system path and entries in the \TeX file database (most \TeX systems use such a database). Thus the “current path” for \TeX is somewhat broader than that for other programs.

For functions which expect a $\langle file\ name \rangle$ argument, this argument may contain both literal items and expandable content, which should on full expansion be the desired file name. Active characters (as declared in `\l_char_active_seq`) are *not* expanded, allowing the direct use of these in file names. Quote tokens (") are not permitted in file names as they are reserved for internal use by some \TeX primitives.

Spaces are trimmed at the beginning and end of the file name: this reflects the fact that some file systems do not allow or interact unpredictably with spaces in these positions. When no extension is given, this will trim spaces from the start of the name only.

1 Input–output stream management

As \TeX engines have a limited number of input and output streams, direct use of the streams by the programmer is not supported in \LaTeX 3. Instead, an internal pool of streams is maintained, and these are allocated and deallocated as needed by other modules. As a result, the programmer should close streams when they are no longer needed, to release them for other processes.

Note that I/O operations are global: streams should all be declared with global names and treated accordingly.

`\ior_new:N`
`\ior_new:c`
`\iow_new:N`
`\iow_new:c`

New: 2011-09-26
Updated: 2011-12-27

`\ior_new:N` $\langle stream \rangle$
`\iow_new:N` $\langle stream \rangle$

Globally reserves the name of the $\langle stream \rangle$, either for reading or for writing as appropriate. The $\langle stream \rangle$ is not opened until the appropriate `\..._open:Nn` function is used. Attempting to use a $\langle stream \rangle$ which has not been opened is an error, and the $\langle stream \rangle$ will behave as the corresponding `\c_term_...`.

`\ior_open:Nn`
`\ior_open:cn`

Updated: 2012-02-10

`\ior_open:Nn` $\langle stream \rangle$ $\{ \langle file\ name \rangle \}$

Opens $\langle file\ name \rangle$ for reading using $\langle stream \rangle$ as the control sequence for file access. If the $\langle stream \rangle$ was already open it is closed before the new operation begins. The $\langle stream \rangle$ is available for access immediately and will remain allocated to $\langle file\ name \rangle$ until a `\ior_close:N` instruction is given or the \TeX run ends. If the file is not found, an error is raised.

<hr/> <code>\ior_open:NnTF</code> <hr/>	<code>\ior_open:NnTF <stream> {<file name>} {<true code>} {<false code>}</code>
<code>\ior_open:cnTF</code> <hr/>	
<hr/> New: 2013-01-12 <hr/>	Opens <i><file name></i> for reading using <i><stream></i> as the control sequence for file access. If the <i><stream></i> was already open it is closed before the new operation begins. The <i><stream></i> is available for access immediately and will remain allocated to <i><file name></i> until a <code>\ior_close:N</code> instruction is given or the TeX run ends. The <i><true code></i> is then inserted into the input stream. If the file is not found, no error is raised and the <i><false code></i> is inserted into the input stream.
<hr/>	
<code>\iow_open:Nn</code> <hr/>	<code>\iow_open:Nn <stream> {<file name>}</code>
<code>\iow_open:cn</code> <hr/>	
<hr/> Updated: 2012-02-09 <hr/>	Opens <i><file name></i> for writing using <i><stream></i> as the control sequence for file access. If the <i><stream></i> was already open it is closed before the new operation begins. The <i><stream></i> is available for access immediately and will remain allocated to <i><file name></i> until a <code>\iow_close:N</code> instruction is given or the TeX run ends. Opening a file for writing clears any existing content in the file (<i>i.e.</i> writing is <i>not</i> additive).
<hr/>	
<code>\ior_close:N</code> <hr/>	<code>\ior_close:N <stream></code>
<code>\ior_close:c</code> <hr/>	<code>\iow_close:N <stream></code>
<code>\iow_close:N</code> <hr/>	Closes the <i><stream></i> . Streams should always be closed when they are finished with as this ensures that they remain available to other programmers.
<code>\iow_close:c</code> <hr/>	
<hr/> Updated: 2012-07-31 <hr/>	
<hr/>	
<code>\ior_show_list:</code> <hr/>	<code>\ior_show_list:</code>
<code>\ior_log_list:</code> <hr/>	<code>\ior_log_list:</code>
<code>\iow_show_list:</code> <hr/>	<code>\iow_show_list:</code>
<code>\iow_log_list:</code> <hr/>	<code>\iow_log_list:</code>
<hr/> New: 2017-06-27 <hr/>	Display (to the terminal or log file) a list of the file names associated with each open (read or write) stream. This is intended for tracking down problems.

1.1 Reading from files

Reading from files and reading from the terminal are separate processes in `expl3`. The functions `\ior_get:NN` and `\ior_str_get:NN`, and their branching equivalents, are designed to work with *files*.

<code>\ior_get:NN</code>
<code>\ior_get:NNTF</code>
New: 2012-06-24
Updated: 2019-03-23

`\ior_get:NN` $\langle stream \rangle$ $\langle token\ list\ variable \rangle$
`\ior_get:NNTF` $\langle stream \rangle$ $\langle token\ list\ variable \rangle$ $\langle true\ code \rangle$ $\langle false\ code \rangle$

Function that reads one or more lines (until an equal number of left and right braces are found) from the file input $\langle stream \rangle$ and stores the result locally in the $\langle token\ list \rangle$ variable. The material read from the $\langle stream \rangle$ is tokenized by \TeX according to the category codes and `\endlinechar` in force when the function is used. Assuming normal settings, any lines which do not end in a comment character `%` have the line ending converted to a space, so for example input

```
a b c
```

results in a token list `a_b_c_`. Any blank line is converted to the token `\par`. Therefore, blank lines can be skipped by using a test such as

```
\ior_get:NN \l_my_stream \l_tmpa_tl
\tl_set:Nn \l_tmpb_tl { \par }
\tl_if_eq:NMF \l_tmpa_tl \l_tmpb_tl
...
```

Also notice that if multiple lines are read to match braces then the resulting token list can contain `\par` tokens. In the non-branching version, where the $\langle stream \rangle$ is not open the $\langle tl\ var \rangle$ is set to `\q_no_value`.

\TeX hackers note: This protected macro is a wrapper around the \TeX primitive `\read`. Regardless of settings, \TeX replaces trailing space and tab characters (character codes 32 and 9) in each line by an end-of-line character (character code `\endlinechar`, omitted if `\endlinechar` is negative or too large) before turning characters into tokens according to current category codes. With default settings, spaces appearing at the beginning of lines are also ignored.

<code>\ior_str_get:NN</code>
<code>\ior_str_get:NNTF</code>
New: 2016-12-04
Updated: 2019-03-23

`\ior_str_get:NN` $\langle stream \rangle$ $\langle token\ list\ variable \rangle$
`\ior_str_get:NNTF` $\langle stream \rangle$ $\langle token\ list\ variable \rangle$ $\langle true\ code \rangle$ $\langle false\ code \rangle$

Function that reads one line from the file input $\langle stream \rangle$ and stores the result locally in the $\langle token\ list \rangle$ variable. The material is read from the $\langle stream \rangle$ as a series of tokens with category code 12 (other), with the exception of space characters which are given category code 10 (space). Multiple whitespace characters are retained by this process. It always only reads one line and any blank lines in the input result in the $\langle token\ list\ variable \rangle$ being empty. Unlike `\ior_get:NN`, line ends do not receive any special treatment. Thus input

```
a b c
```

results in a token list `a b c` with the letters `a`, `b`, and `c` having category code 12. In the non-branching version, where the $\langle stream \rangle$ is not open the $\langle tl\ var \rangle$ is set to `\q_no_value`.

\TeX hackers note: This protected macro is a wrapper around the $\varepsilon\text{\TeX}$ primitive `\readline`. Regardless of settings, \TeX removes trailing space and tab characters (character codes 32 and 9). However, the end-line character normally added by this primitive is not included in the result of `\ior_str_get:NN`.

All mappings are done at the current group level, *i.e.* any local assignments made by the $\langle function \rangle$ or $\langle code \rangle$ discussed below remain in effect after the loop.

<hr/> <code>\ior_map_inline:Nn</code> <hr/>	<code>\ior_map_inline:Nn <stream> {<inline function>}</code>
<hr/> New: 2012-02-11 <hr/>	Applies the <i><inline function></i> to each set of <i><lines></i> obtained by calling <code>\ior_get:NN</code> until reaching the end of the file. $\mathrm{T\!E\!X}$ ignores any trailing new-line marker from the file it reads. The <i><inline function></i> should consist of code which receives the <i><line></i> as <i>#1</i> .
<hr/> <code>\ior_str_map_inline:Nn</code> <hr/>	<code>\ior_str_map_inline:Nn <stream> {<inline function>}</code>
<hr/> New: 2012-02-11 <hr/>	Applies the <i><inline function></i> to every <i><line></i> in the <i><stream></i> . The material is read from the <i><stream></i> as a series of tokens with category code 12 (other), with the exception of space characters which are given category code 10 (space). The <i><inline function></i> should consist of code which receives the <i><line></i> as <i>#1</i> . Note that $\mathrm{T\!E\!X}$ removes trailing space and tab characters (character codes 32 and 9) from every line upon input. $\mathrm{T\!E\!X}$ also ignores any trailing new-line marker from the file it reads.
<hr/> <code>\ior_map_variable:NNn</code> <hr/>	<code>\ior_map_variable:NNn <stream> <tl var> {<code>}</code>
<hr/> New: 2019-01-13 <hr/>	For each set of <i><lines></i> obtained by calling <code>\ior_get:NN</code> until reaching the end of the file, stores the <i><lines></i> in the <i><tl var></i> then applies the <i><code></i> . The <i><code></i> will usually make use of the <i><variable></i> , but this is not enforced. The assignments to the <i><variable></i> are local. Its value after the loop is the last set of <i><lines></i> , or its original value if the <i><stream></i> is empty. $\mathrm{T\!E\!X}$ ignores any trailing new-line marker from the file it reads. This function is typically faster than <code>\ior_map_inline:Nn</code> .
<hr/> <code>\ior_str_map_variable:NNn</code> <hr/>	<code>\ior_str_map_variable:NNn <stream> <variable> {<code>}</code>
<hr/> New: 2019-01-13 <hr/>	For each <i><line></i> in the <i><stream></i> , stores the <i><line></i> in the <i><variable></i> then applies the <i><code></i> . The material is read from the <i><stream></i> as a series of tokens with category code 12 (other), with the exception of space characters which are given category code 10 (space). The <i><code></i> will usually make use of the <i><variable></i> , but this is not enforced. The assignments to the <i><variable></i> are local. Its value after the loop is the last <i><line></i> , or its original value if the <i><stream></i> is empty. Note that $\mathrm{T\!E\!X}$ removes trailing space and tab characters (character codes 32 and 9) from every line upon input. $\mathrm{T\!E\!X}$ also ignores any trailing new-line marker from the file it reads. This function is typically faster than <code>\ior_str_map_inline:Nn</code> .
<hr/> <code>\ior_map_break:</code> <hr/>	<code>\ior_map_break:</code>
<hr/> New: 2012-06-29 <hr/>	Used to terminate a <code>\ior_map_...</code> function before all lines from the <i><stream></i> have been processed. This normally takes place within a conditional statement, for example <pre> \ior_map_inline:Nn \l_my_ior { \str_if_eq:nnTF { #1 } { bingo } { \ior_map_break: } { % Do something useful } }</pre>

Use outside of a `\ior_map_...` scenario leads to low level $\mathrm{T\!E\!X}$ errors.

$\mathrm{T\!E\!X}$ hackers note: When the mapping is broken, additional tokens may be inserted before further items are taken from the input stream. This depends on the design of the mapping function.

`\ior_map_break:n`

`New: 2012-06-29`

`\ior_map_break:n {<code>}`

Used to terminate a `\ior_map_...` function before all lines in the `<stream>` have been processed, inserting the `<code>` after the mapping has ended. This normally takes place within a conditional statement, for example

```
\ior_map_inline:Nn \l_my_ior
{
  \str_if_eq:nnTF { #1 } { bingo }
  { \ior_map_break:n { <code> } }
  {
    % Do something useful
  }
}
```

Use outside of a `\ior_map_...` scenario leads to low level `TeX` errors.

TeXhackers note: When the mapping is broken, additional tokens may be inserted before the `<code>` is inserted into the input stream. This depends on the design of the mapping function.

`\ior_if_eof_p:N *``\ior_if_eof:NTF *`

`Updated: 2012-02-10`

`\ior_if_eof_p:N <stream>``\ior_if_eof:NTF <stream> {<true code>} {<false code>}`

Tests if the end of a file `<stream>` has been reached during a reading operation. The test also returns a `true` value if the `<stream>` is not open.

1.2 Writing to files

`\iow_now:Nn``\iow_now:(Nx|cn|cx)`

`Updated: 2012-06-05`

`\iow_now:Nn <stream> {<tokens>}`

This functions writes `<tokens>` to the specified `<stream>` immediately (*i.e.* the write operation is called on expansion of `\iow_now:Nn`).

`\iow_log:n``\iow_log:x``\iow_log:n {<tokens>}`

This function writes the given `<tokens>` to the log (transcript) file immediately: it is a dedicated version of `\iow_now:Nn`.

`\iow_term:n``\iow_term:x``\iow_term:n {<tokens>}`

This function writes the given `<tokens>` to the terminal file immediately: it is a dedicated version of `\iow_now:Nn`.

<hr/>	
<code>\iow_shipout:Nn</code>	<code>\iow_shipout:Nn <stream> {\tokens}</code>
<code>\iow_shipout:(Nx cn cx)</code>	This functions writes $\langle tokens \rangle$ to the specified $\langle stream \rangle$ when the current page is finalised (<i>i.e.</i> at shipout). The x -type variants expand the $\langle tokens \rangle$ at the point where the function is used but <i>not</i> when the resulting tokens are written to the $\langle stream \rangle$ (<i>cf.</i> <code>\iow_shipout_x:Nn</code>).
<hr/>	

T_EXhackers note: When using `expl3` with a format other than L^AT_EX, new line characters inserted using `\iow_newline:` or using the line-wrapping code `\iow_wrap:nnnN` are not recognized in the argument of `\iow_shipout:Nn`. This may lead to the insertion of additional unwanted line-breaks.

<hr/>	
<code>\iow_shipout_x:Nn</code>	<code>\iow_shipout_x:Nn <stream> {\tokens}</code>
<code>\iow_shipout_x:(Nx cn cx)</code>	This functions writes $\langle tokens \rangle$ to the specified $\langle stream \rangle$ when the current page is finalised (<i>i.e.</i> at shipout). The $\langle tokens \rangle$ are expanded at the time of writing in addition to any expansion when the function is used. This makes these functions suitable for including material finalised during the page building process (such as the page number integer).
Updated: 2012-09-08	
<hr/>	

T_EXhackers note: This is a wrapper around the T_EX primitive `\write`. When using `expl3` with a format other than L^AT_EX, new line characters inserted using `\iow_newline:` or using the line-wrapping code `\iow_wrap:nnnN` are not recognized in the argument of `\iow_shipout:Nn`. This may lead to the insertion of additional unwanted line-breaks.

<hr/>	
<code>\iow_char:N *</code>	<code>\iow_char:N \<char></code>
	Inserts $\langle char \rangle$ into the output stream. Useful when trying to write difficult characters such as %, {, }, <i>etc.</i> in messages, for example:

`\iow_now:Nx \g_my_iow { \iow_char:N \{ text \iow_char:N \} }`

The function has no effect if writing is taking place without expansion (*e.g.* in the second argument of `\iow_now:Nn`).

<hr/>	
<code>\iow_newline: *</code>	<code>\iow_newline:</code>
	Function to add a new line within the $\langle tokens \rangle$ written to a file. The function has no effect if writing is taking place without expansion (<i>e.g.</i> in the second argument of <code>\iow_now:Nn</code>).

T_EXhackers note: When using `expl3` with a format other than L^AT_EX, the character inserted by `\iow_newline:` is not recognized by T_EX, which may lead to the insertion of additional unwanted line-breaks. This issue only affects `\iow_shipout:Nn`, `\iow_shipout_x:Nn` and direct uses of primitive operations.

1.3 Wrapping lines in output

`\iow_wrap:nnnN`
`\iow_wrap:nxnN`

New: 2012-06-28
Updated: 2017-12-04

`\iow_wrap:nnnN` $\langle text \rangle$ $\langle run-on\ text \rangle$ $\langle set\ up \rangle$ $\langle function \rangle$

This function wraps the $\langle text \rangle$ to a fixed number of characters per line. At the start of each line which is wrapped, the $\langle run-on\ text \rangle$ is inserted. The line character count targeted is the value of `\l_iow_line_count_int` minus the number of characters in the $\langle run-on\ text \rangle$ for all lines except the first, for which the target number of characters is simply `\l_iow_line_count_int` since there is no run-on text. The $\langle text \rangle$ and $\langle run-on\ text \rangle$ are exhaustively expanded by the function, with the following substitutions:

- `\` or `\iow_newline`: may be used to force a new line,
- `_` may be used to represent a forced space (for example after a control sequence),
- `\#`, `\%`, `\{`, `\}`, `\~` may be used to represent the corresponding character,
- `\iow_allow_break`: may be used to allow a line-break without inserting a space (this is experimental),
- `\iow_indent:n` may be used to indent a part of the $\langle text \rangle$ (not the $\langle run-on\ text \rangle$).

Additional functions may be added to the wrapping by using the $\langle set\ up \rangle$, which is executed before the wrapping takes place: this may include overriding the substitutions listed.

Any expandable material in the $\langle text \rangle$ which is not to be expanded on wrapping should be converted to a string using `\token_to_str:N`, `\tl_to_str:n`, `\tl_to_str:N`, etc.

The result of the wrapping operation is passed as a braced argument to the $\langle function \rangle$, which is typically a wrapper around a write operation. The output of `\iow_wrap:nnnN` (i.e. the argument passed to the $\langle function \rangle$) consists of characters of category “other” (category code 12), with the exception of spaces which have category “space” (category code 10). This means that the output does *not* expand further when written to a file.

T_EXhackers note: Internally, `\iow_wrap:nnnN` carries out an x-type expansion on the $\langle text \rangle$ to expand it. This is done in such a way that `\exp_not:N` or `\exp_not:n` could be used to prevent expansion of material. However, this is less conceptually clear than conversion to a string, which is therefore the supported method for handling expandable material in the $\langle text \rangle$.

`\iow_indent:n`

New: 2011-09-21

`\iow_indent:n` $\langle text \rangle$

In the first argument of `\iow_wrap:nnnN` (for instance in messages), indents $\langle text \rangle$ by four spaces. This function does not cause a line break, and only affects lines which start within the scope of the $\langle text \rangle$. In case the indented $\langle text \rangle$ should appear on separate lines from the surrounding text, use `\` to force line breaks.

`\l_iow_line_count_int`

New: 2012-06-24

The maximum number of characters in a line to be written by the `\iow_wrap:nnnN` function. This value depends on the T_EX system in use: the standard value is 78, which is typically correct for unmodified T_EX Live and MiK_T_EX systems.

1.4 Constant input–output streams, and variables

<code>\g_tmpa_iow</code>	Scratch input stream for global use. These are never used by the kernel code, and so are safe for use with any L ^A T _E X3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.
<code>\g_tmpb_iow</code>	
<hr/> New: 2017-12-11 <hr/>	

<code>\c_log_iow</code>	Constant output streams for writing to the log and to the terminal (plus the log), respectively.
<code>\c_term_iow</code>	

<code>\g_tmpa_iow</code>	Scratch output stream for global use. These are never used by the kernel code, and so are safe for use with any L ^A T _E X3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.
<code>\g_tmpb_iow</code>	
<hr/> New: 2017-12-11 <hr/>	

1.5 Primitive conditionals

<code>\if_eof:w ★</code>	<code>\if_eof:w <stream></code> <code> <true code></code> <code>\else:</code> <code> <false code></code> <code>\fi:</code>
Tests if the <code><stream></code> returns “end of file”, which is true for non-existent files. The <code>\else:</code> branch is optional.	

T_EXhackers note: This is the T_EX primitive `\ifeof`.

2 File operation functions

<code>\g_file_curr_dir_str</code> <code>\g_file_curr_name_str</code> <code>\g_file_curr_ext_str</code>	Contain the directory, name and extension of the current file. The directory is empty if the file was loaded without an explicit path (<i>i.e.</i> if it is in the T _E X search path), and does <i>not</i> end in / other than the case that it is exactly equal to the root directory. The <code><name></code> and <code><ext></code> parts together make up the file name, thus the <code><name></code> part may be thought of as the “job name” for the current file. Note that T _E X does not provide information on the <code><ext></code> part for the main (top level) file and that this file always has an empty <code><dir></code> component. Also, the <code><name></code> here will be equal to <code>\c_sys_jobname_str</code> , which may be different from the real file name (if set using <code>--jobname</code> , for example).
<hr/> New: 2017-06-21 <hr/>	

<hr/> <code>\l_file_search_path_seq</code> <hr/> New: 2017-06-18	<p>Each entry is the path to a directory which should be searched when seeking a file. Each path can be relative or absolute, and should not include the trailing slash. The entries are not expanded when used so may contain active characters but should not feature any variable content. Spaces need not be quoted.</p> <p>T_EXhackers note: When working as a package in L^AT_EX 2_ε, <code>expl3</code> will automatically append the current <code>\input@path</code> to the set of values from <code>\l_file_search_path_seq</code>.</p>
<hr/> <code>\file_if_exist:nTF</code> <hr/> Updated: 2012-02-10	<code>\file_if_exist:nTF {<file name>} {<true code>} {<false code>}</code> <p>Searches for <code><file name></code> using the current T_EX search path and the additional paths controlled by <code>\l_file_search_path_seq</code>.</p>
<hr/> <code>\file_get:nnN</code> <code>\file_get:nnNTF</code> <hr/> New: 2019-01-16 Updated: 2019-02-16	<code>\file_get:nnN {<filename>} {<setup>} <tl></code> <code>\file_get:nnNTF {<filename>} {<setup>} <tl> {<true code>} {<false code>}</code> <p>Defines <code><tl></code> to the contents of <code><filename></code>. Category codes may need to be set appropriately via the <code><setup></code> argument. The non-branching version sets the <code><tl></code> to <code>\q_no_value</code> if the file is not found. The branching version runs the <code><true code></code> after the assignment to <code><tl></code> if the file is found, and <code><false code></code> otherwise.</p>
<hr/> <code>\file_get_full_name:nN</code> <code>\file_get_full_name:VN</code> <code>\file_get_full_name:nNTF</code> <code>\file_get_full_name:VNTF</code> <hr/> Updated: 2019-02-16	<code>\file_get_full_name:nN {<file name>} <tl></code> <code>\file_get_full_name:nNTF {<file name>} <tl> {<true code>} {<false code>}</code> <p>Searches for <code><file name></code> in the path as detailed for <code>\file_if_exist:nTF</code>, and if found sets the <code><tl var></code> the fully-qualified name of the file, <i>i.e.</i> the path and file name. This includes an extension <code>.tex</code> when the given <code><file name></code> has no extension but the file found has that extension. In the non-branching version, the <code><tl var></code> will be set to <code>\q_no_value</code> in the case that the file does not exist.</p>
<hr/> <code>\file_full_name:n</code> ☆ <code>\file_full_name:V</code> ☆ <hr/> New: 2019-09-03	<code>\file_full_name:n {<file name>}</code> <p>Searches for <code><file name></code> in the path as detailed for <code>\file_if_exist:nTF</code>, and if found leaves the fully-qualified name of the file, <i>i.e.</i> the path and file name, in the input stream. This includes an extension <code>.tex</code> when the given <code><file name></code> has no extension but the file found has that extension. If the file is not found on the path, the expansion is empty.</p>

`\file_parse_full_name:nNNN`
`\file_parse_full_name:VNNN`

New: 2017-06-23
Updated: 2020-06-24

`\file_parse_full_name:nNNN {<full name>} <dir> <name> <ext>`

Parses the `<full name>` and splits it into three parts, each of which is returned by setting the appropriate local string variable:

- The `<dir>`: everything up to the last / (path separator) in the `<file path>`. As with system PATH variables and related functions, the `<dir>` does *not* include the trailing / unless it points to the root directory. If there is no path (only a file name), `<dir>` is empty.
- The `<name>`: everything after the last / up to the last ., where both of those characters are optional. The `<name>` may contain multiple . characters. It is empty if `<full name>` consists only of a directory name.
- The `<ext>`: everything after the last . (including the dot). The `<ext>` is empty if there is no . after the last /.

Before parsing, the `<full name>` is expanded until only non-expandable tokens remain, except that active characters are also not expanded. Quotes (") are invalid in file names and are discarded from the input.

`\file_parse_full_name:n *`

New: 2020-06-24

`\file_parse_full_name:n {<full name>}`

Parses the `<full name>` as described for `\file_parse_full_name:nNNN`, and leaves `<dir>`, `<name>`, and `<ext>` in the input stream, each inside a pair of braces.

`\file_parse_full_name_apply:nN *` `\file_parse_full_name_apply:nN {<full name>} <function>`

New: 2020-06-24

Parses the `<full name>` as described for `\file_parse_full_name:nNNN`, and passes `<dir>`, `<name>`, and `<ext>` as arguments to `<function>`, as an n-type argument each, in this order.

`\file_hex_dump:n` ☆
`\file_hex_dump:nnn` ☆

New: 2019-11-19

`\file_hex_dump:n {<file name>}`
`\file_hex_dump:nnn {<file name>} {<start index>} {<end index>}`

Searches for `<file name>` using the current T_EX search path and the additional paths controlled by `\l_file_search_path_seq`. It then expands to leave the hexadecimal dump of the file content in the input stream. The file is read as bytes, which means that in contrast to most T_EX behaviour there will be a difference in result depending on the line endings used in text files. The same file will produce the same result between different engines: the algorithm used is the same in all cases. When the file is not found, the result of expansion is empty. The `{<start index>}` and `{<end index>}` values work as described for `\str_range:nnn`.

`\file_get_hex_dump:nN`
`\file_get_hex_dump:nNTF`
`\file_get_hex_dump:nnnN`
`\file_get_hex_dump:nnnNTF`

New: 2019-11-19

`\file_get_hex_dump:nN {<file name>} <tl var>`
`\file_get_hex_dump:nnnN {<file name>} {<start index>} {<end index>} <tl var>`

Sets the `<tl var>` to the result of applying `\file_hex_dump:n/\file_hex_dump:nnn` to the `<file>`. If the file is not found, the `<tl var>` will be set to `\q_no_value`.

<hr/> <code>\file_md5hash:n</code> ☆	<code>\file_md5hash:n {⟨file name⟩}</code>
<hr/> New: 2019-09-03	<p>Searches for $\langle file\ name \rangle$ using the current \TeX search path and the additional paths controlled by <code>\l_file_search_path_seq</code>. It then expands to leave the MD5 sum generated from the contents of the file in the input stream. The file is read as bytes, which means that in contrast to most \TeX behaviour there will be a difference in result depending on the line endings used in text files. The same file will produce the same result between different engines: the algorithm used is the same in all cases. When the file is not found, the result of expansion is empty.</p>
<hr/> <code>\file_get_md5hash:nN</code> <code>\file_get_md5hash:nNTF</code>	<code>\file_get_md5hash:nN {⟨file name⟩} ⟨tl var⟩</code>
<hr/> New: 2017-07-11 Updated: 2019-02-16	<p>Sets the $\langle tl\ var \rangle$ to the result of applying <code>\file_md5hash:n</code> to the $\langle file \rangle$. If the file is not found, the $\langle tl\ var \rangle$ will be set to <code>\q_no_value</code>.</p>
<hr/> <code>\file_size:n</code> ☆	<code>\file_size:n {⟨file name⟩}</code>
<hr/> New: 2019-09-03	<p>Searches for $\langle file\ name \rangle$ using the current \TeX search path and the additional paths controlled by <code>\l_file_search_path_seq</code>. It then expands to leave the size of the file in bytes in the input stream. When the file is not found, the result of expansion is empty.</p>
<hr/> <code>\file_get_size:nN</code> <code>\file_get_size:nNTF</code>	<code>\file_get_size:nN {⟨file name⟩} ⟨tl var⟩</code>
<hr/> New: 2017-07-09 Updated: 2019-02-16	<p>Sets the $\langle tl\ var \rangle$ to the result of applying <code>\file_size:n</code> to the $\langle file \rangle$. If the file is not found, the $\langle tl\ var \rangle$ will be set to <code>\q_no_value</code>. This is not available in older versions of \LaTeX.</p>
<hr/> <code>\file_timestamp:n</code> ☆	<code>\file_timestamp:n {⟨file name⟩}</code>
<hr/> New: 2019-09-03	<p>Searches for $\langle file\ name \rangle$ using the current \TeX search path and the additional paths controlled by <code>\l_file_search_path_seq</code>. It then expands to leave the modification timestamp of the file in the input stream. The timestamp is of the form $D:\langle year \rangle \langle month \rangle \langle day \rangle \langle hour \rangle \langle minute \rangle \langle second \rangle \langle offset \rangle$, where the latter may be Z (UTC) or $\langle plus-minus \rangle \langle hours \rangle' \langle minutes \rangle'$. When the file is not found, the result of expansion is empty. This is not available in older versions of \LaTeX.</p>
<hr/> <code>\file_get_timestamp:nN</code> <code>\file_get_timestamp:nNTF</code>	<code>\file_get_timestamp:nN {⟨file name⟩} ⟨tl var⟩</code>
<hr/> New: 2017-07-09 Updated: 2019-02-16	<p>Sets the $\langle tl\ var \rangle$ to the result of applying <code>\file_timestamp:n</code> to the $\langle file \rangle$. If the file is not found, the $\langle tl\ var \rangle$ will be set to <code>\q_no_value</code>. This is not available in older versions of \LaTeX.</p>

<code>\file_compare_timestamp_p:nNn</code>	<code>★</code>	<code>\file_compare_timestamp:nNn {<file-1>} <comparator> {<file-2>} {<true</code>
<code>\file_compare_timestamp:nNnTF</code>	<code>★</code>	<code>code)} {<false code>}</code>

New: 2019-05-13

Updated: 2019-09-20

Compares the file stamps on the two *<files>* as indicated by the *<comparator>*, and inserts either the *<true code>* or *<false case>* as required. A file which is not found is treated as older than any file which is found. This allows for example the construct

```
\file_compare_timestamp:nNnT { source-file } > { derived-file }
{
  % Code to regenerate derived file
}
```

to work when the derived file is entirely absent. The timestamp of two absent files is regarded as different. This is not available in older versions of X_YTeX.

<code>\file_input:n</code>	<code>\file_input:n {<file name>}</code>
----------------------------	--

Updated: 2017-06-26

Searches for *<file name>* in the path as detailed for `\file_if_exist:nTF`, and if found reads in the file as additional L^AT_EX source. All files read are recorded for information and the file name stack is updated by this function. An error is raised if the file is not found.

<code>\file_if_exist_input:n</code>	<code>\file_if_exist_input:n {<file name>}</code>
<code>\file_if_exist_input:nF</code>	<code>\file_if_exist_input:nF {<file name>} {<false code>}</code>

New: 2014-07-02

Searches for *<file name>* using the current T_EX search path and the additional paths controlled by `\file_path_include:n`. If found then reads in the file as additional L^AT_EX source as described for `\file_input:n`, otherwise inserts the *<false code>*. Note that these functions do not raise an error if the file is not found, in contrast to `\file_input:n`.

<code>\file_input_stop:</code>	<code>\file_input_stop:</code>
--------------------------------	--------------------------------

New: 2017-07-07

Ends the reading of a file started by `\file_input:n` or similar before the end of the file is reached. Where the file reading is being terminated due to an error, `\msg_critical:nn(nn)` should be preferred.

T_EXhackers note: This function must be used on a line on its own: T_EX reads files line-by-line and so any additional tokens in the “current” line will still be read.

This is also true if the function is hidden inside another function (which will be the normal case), i.e., all tokens on the same line in the source file are still processed. Putting it on a line by itself in the definition doesn’t help as it is the line where it is used that counts!

<code>\file_show_list:</code>	<code>\file_show_list:</code>
<code>\file_log_list:</code>	<code>\file_log_list:</code>

These functions list all files loaded by L^AT_EX 2_ε commands that populate `\@filelist` or by `\file_input:n`. While `\file_show_list:` displays the list in the terminal, `\file_log_list:` outputs it to the log file only.

Part XX

The l3skip package

Dimensions and skips

L^AT_EX3 provides two general length variables: `dim` and `skip`. Lengths stored as `dim` variables have a fixed length, whereas `skip` lengths have a rubber (stretch/shrink) component. In addition, the `muskip` type is available for use in math mode: this is a special form of `skip` where the lengths involved are determined by the current math font (in μ). There are common features in the creation and setting of length variables, but for clarity the functions are grouped by variable type.

1 Creating and initialising `dim` variables

<code>\dim_new:N</code>
<code>\dim_new:c</code>

`\dim_new:N` $\langle dimension \rangle$

Creates a new $\langle dimension \rangle$ or raises an error if the name is already taken. The declaration is global. The $\langle dimension \rangle$ is initially equal to 0pt.

<code>\dim_const:Nn</code>
<code>\dim_const:cn</code>

`\dim_const:Nn` $\langle dimension \rangle$ $\{ \langle dimension expression \rangle \}$

Creates a new constant $\langle dimension \rangle$ or raises an error if the name is already taken. The value of the $\langle dimension \rangle$ is set globally to the $\langle dimension expression \rangle$.

New: 2012-03-05

<code>\dim_zero:N</code>
<code>\dim_zero:c</code>
<code>\dim_gzero:N</code>
<code>\dim_gzero:c</code>

`\dim_zero:N` $\langle dimension \rangle$

Sets $\langle dimension \rangle$ to 0pt.

<code>\dim_zero_new:N</code>
<code>\dim_zero_new:c</code>
<code>\dim_gzero_new:N</code>
<code>\dim_gzero_new:c</code>

`\dim_zero_new:N` $\langle dimension \rangle$

Ensures that the $\langle dimension \rangle$ exists globally by applying `\dim_new:N` if necessary, then applies `\dim_(g)zero:N` to leave the $\langle dimension \rangle$ set to zero.

New: 2012-01-07

<code>\dim_if_exist_p:N</code> \star
<code>\dim_if_exist_p:c</code> \star
<code>\dim_if_exist:N\overline{TF}</code> \star
<code>\dim_if_exist:c\overline{TF}</code> \star

`\dim_if_exist_p:N` $\langle dimension \rangle$

`\dim_if_exist:NTF` $\langle dimension \rangle$ $\{ \langle true code \rangle \} \{ \langle false code \rangle \}$

Tests whether the $\langle dimension \rangle$ is currently defined. This does not check that the $\langle dimension \rangle$ really is a dimension variable.

New: 2012-03-03

2 Setting dim variables

<code>\dim_add:Nn</code>	<code>\dim_add:Nn <dimension> {<dimension expression>}</code>
<code>\dim_add:cn</code>	
<code>\dim_gadd:Nn</code>	Adds the result of the $\langle dimension\ expression \rangle$ to the current content of the $\langle dimension \rangle$.
<code>\dim_gadd:cn</code>	

Updated: 2011-10-22

<code>\dim_set:Nn</code>	<code>\dim_set:Nn <dimension> {<dimension expression>}</code>
<code>\dim_set:cn</code>	
<code>\dim_gset:Nn</code>	Sets $\langle dimension \rangle$ to the value of $\langle dimension\ expression \rangle$, which must evaluate to a length with units.
<code>\dim_gset:cn</code>	

Updated: 2011-10-22

<code>\dim_set_eq:NN</code>	<code>\dim_set_eq:NN <dimension₁₂</code>
<code>\dim_set_eq:(cN Nc cc)</code>	Sets the content of $\langle dimension_1 \rangle$ equal to that of $\langle dimension_2 \rangle$.
<code>\dim_gset_eq:NN</code>	
<code>\dim_gset_eq:(cN Nc cc)</code>	

<code>\dim_sub:Nn</code>	<code>\dim_sub:Nn <dimension> {<dimension expression>}</code>
<code>\dim_sub:cn</code>	
<code>\dim_gsub:Nn</code>	Subtracts the result of the $\langle dimension\ expression \rangle$ from the current content of the $\langle dimension \rangle$.
<code>\dim_gsub:cn</code>	

Updated: 2011-10-22

3 Utilities for dimension calculations

<code>\dim_abs:n</code>	<code>\dim_abs:n {<dimexpr>}</code>
Updated: 2012-09-26	Converts the $\langle dimexpr \rangle$ to its absolute value, leaving the result in the input stream as a $\langle dimension\ denotation \rangle$.

<code>\dim_max:nn</code>	<code>\dim_max:nn {<dimexpr₁>} {<dimexpr₂>}</code>
<code>\dim_min:nn</code>	<code>\dim_min:nn {<dimexpr₁>} {<dimexpr₂>}</code>
New: 2012-09-09	Evaluates the two $\langle dimension\ expressions \rangle$ and leaves either the maximum or minimum value in the input stream as appropriate, as a $\langle dimension\ denotation \rangle$.
Updated: 2012-09-26	

`\dim_ratio:nn` ☆

Updated: 2011-10-22

`\dim_ratio:nn` { $\langle dimexpr_1 \rangle$ } { $\langle dimexpr_2 \rangle$ }

Parses the two $\langle dimension expressions \rangle$ and converts the ratio of the two to a form suitable for use inside a $\langle dimension expression \rangle$. This ratio is then left in the input stream, allowing syntax such as

```
\dim_set:Nn \l_my_dim
{ 10 pt * \dim_ratio:nn { 5 pt } { 10 pt } }
```

The output of `\dim_ratio:nn` on full expansion is a ratio expression between two integers, with all distances converted to scaled points. Thus

```
\tl_set:Nx \l_my_tl { \dim_ratio:nn { 5 pt } { 10 pt } }
\tl_show:N \l_my_tl
```

displays 327680/655360 on the terminal.

4 Dimension expression conditionals

`\dim_compare_p:nNn` ★

`\dim_compare:nNnTF` ★

`\dim_compare_p:nNn` { $\langle dimexpr_1 \rangle$ } $\langle relation \rangle$ { $\langle dimexpr_2 \rangle$ }

`\dim_compare:nNnTF`

{ $\langle dimexpr_1 \rangle$ } $\langle relation \rangle$ { $\langle dimexpr_2 \rangle$ }
{ $\langle true code \rangle$ } { $\langle false code \rangle$ }

This function first evaluates each of the $\langle dimension expressions \rangle$ as described for `\dim_eval:n`. The two results are then compared using the $\langle relation \rangle$:

Equal	=
Greater than	>
Less than	<

This function is less flexible than `\dim_compare:nTF` but around 5 times faster.

```

\dim_compare_p:n * \dim_compare_p:n
\dim_compare:nTF * {
    <dimexpr1> <relation1>
    ...
    <dimexprN> <relationN>
    <dimexprN+1>
}
\dim_compare:nTF
{
    <dimexpr1> <relation1>
    ...
    <dimexprN> <relationN>
    <dimexprN+1>
}
{{true code}} {{false code}}

```

Updated: 2013-01-13

This function evaluates the *<dimension expressions>* as described for `\dim_eval:n` and compares consecutive result using the corresponding *<relation>*, namely it compares *<dimexpr₁>* and *<dimexpr₂>* using the *<relation₁>*, then *<dimexpr₂>* and *<dimexpr₃>* using the *<relation₂>*, until finally comparing *<dimexpr_N>* and *<dimexpr_{N+1}>* using the *<relation_N>*. The test yields `true` if all comparisons are `true`. Each *<dimension expression>* is evaluated only once, and the evaluation is lazy, in the sense that if one comparison is `false`, then no other *<dimension expression>* is evaluated and no other comparison is performed. The *<relations>* can be any of the following:

Equal	= or ==
Greater than or equal to	>=
Greater than	>
Less than or equal to	<=
Less than	<
Not equal	!=

This function is more flexible than `\dim_compare:nNnTF` but around 5 times slower.

<code>\dim_case:nn</code> ☆	<code>\dim_case:nnTF {⟨test dimension expression⟩}</code>
<code>\dim_case:nnTF</code> ☆	<code>{</code>
New: 2013-07-24	<code>{⟨dimexpr case₁⟩} {⟨code case₁⟩}</code>
	<code>{⟨dimexpr case₂⟩} {⟨code case₂⟩}</code>
	<code>...</code>
	<code>{⟨dimexpr case_n⟩} {⟨code case_n⟩}</code>
	<code>}</code>
	<code>{⟨true code⟩}</code>
	<code>{⟨false code⟩}</code>

This function evaluates the *⟨test dimension expression⟩* and compares this in turn to each of the *⟨dimension expression cases⟩*. If the two are equal then the associated *⟨code⟩* is left in the input stream and other cases are discarded. If any of the cases are matched, the *⟨true code⟩* is also inserted into the input stream (after the code for the appropriate case), while if none match then the *⟨false code⟩* is inserted. The function `\dim_case:nn`, which does nothing if there is no match, is also available. For example

```

\dim_set:Nn \l_tmpa_dim { 5 pt }
\dim_case:nnF
{ 2 \l_tmpa_dim }
{
  { 5 pt }      { Small }
  { 4 pt + 6 pt } { Medium }
  { - 10 pt }   { Negative }
}
{ No idea! }
```

leaves “Medium” in the input stream.

5 Dimension expression loops

<code>\dim_do_until:nNnn</code> ☆	<code>\dim_do_until:nNnn {⟨dimexpr₁⟩} ⟨relation⟩ {⟨dimexpr₂⟩} {⟨code⟩}</code>
-----------------------------------	---

Places the *⟨code⟩* in the input stream for T_EX to process, and then evaluates the relationship between the two *⟨dimension expressions⟩* as described for `\dim_compare:nNnTF`. If the test is **false** then the *⟨code⟩* is inserted into the input stream again and a loop occurs until the *⟨relation⟩* is **true**.

<code>\dim_do_while:nNnn</code> ☆	<code>\dim_do_while:nNnn {⟨dimexpr₁⟩} ⟨relation⟩ {⟨dimexpr₂⟩} {⟨code⟩}</code>
-----------------------------------	---

Places the *⟨code⟩* in the input stream for T_EX to process, and then evaluates the relationship between the two *⟨dimension expressions⟩* as described for `\dim_compare:nNnTF`. If the test is **true** then the *⟨code⟩* is inserted into the input stream again and a loop occurs until the *⟨relation⟩* is **false**.

<code>\dim_until_do:nNnn</code> ☆	<code>\dim_until_do:nNnn {⟨dimexpr₁⟩} ⟨relation⟩ {⟨dimexpr₂⟩} {⟨code⟩}</code>
-----------------------------------	---

Evaluates the relationship between the two *⟨dimension expressions⟩* as described for `\dim_compare:nNnTF`, and then places the *⟨code⟩* in the input stream if the *⟨relation⟩* is **false**. After the *⟨code⟩* has been processed by T_EX the test is repeated, and a loop occurs until the test is **true**.

<hr/> <code>\dim_while_do:nNnn</code> ☆ <hr/>	<code>\dim_while_do:nNnn {<dimexpr₁>} <relation> {<dimexpr₂>} {<code>}</code>
	Evaluates the relationship between the two <i><dimension expressions></i> as described for <code>\dim_compare:nNnTF</code> , and then places the <i><code></i> in the input stream if the <i><relation></i> is true . After the <i><code></i> has been processed by T _E X the test is repeated, and a loop occurs until the test is false .
<hr/> <code>\dim_do_until:nn</code> ☆ <hr/> <div>Updated: 2013-01-13</div>	<code>\dim_do_until:nn {<dimension relation>} {<code>}</code>
	Places the <i><code></i> in the input stream for T _E X to process, and then evaluates the <i><dimension relation></i> as described for <code>\dim_compare:nTF</code> . If the test is false then the <i><code></i> is inserted into the input stream again and a loop occurs until the <i><relation></i> is true .
<hr/> <code>\dim_do_while:nn</code> ☆ <hr/> <div>Updated: 2013-01-13</div>	<code>\dim_do_while:nn {<dimension relation>} {<code>}</code>
	Places the <i><code></i> in the input stream for T _E X to process, and then evaluates the <i><dimension relation></i> as described for <code>\dim_compare:nTF</code> . If the test is true then the <i><code></i> is inserted into the input stream again and a loop occurs until the <i><relation></i> is false .
<hr/> <code>\dim_until_do:nn</code> ☆ <hr/> <div>Updated: 2013-01-13</div>	<code>\dim_until_do:nn {<dimension relation>} {<code>}</code>
	Evaluates the <i><dimension relation></i> as described for <code>\dim_compare:nTF</code> , and then places the <i><code></i> in the input stream if the <i><relation></i> is false . After the <i><code></i> has been processed by T _E X the test is repeated, and a loop occurs until the test is true .
<hr/> <code>\dim_while_do:nn</code> ☆ <hr/> <div>Updated: 2013-01-13</div>	<code>\dim_while_do:nn {<dimension relation>} {<code>}</code>
	Evaluates the <i><dimension relation></i> as described for <code>\dim_compare:nTF</code> , and then places the <i><code></i> in the input stream if the <i><relation></i> is true . After the <i><code></i> has been processed by T _E X the test is repeated, and a loop occurs until the test is false .

6 Dimension step functions

<hr/> <code>\dim_step_function:nnnN</code> ☆ <hr/> <div>New: 2018-02-18</div>	<code>\dim_step_function:nnnN {<initial value>} {<step>} {<final value>} <function></code>
	This function first evaluates the <i><initial value></i> , <i><step></i> and <i><final value></i> , all of which should be dimension expressions. The <i><function></i> is then placed in front of each <i><value></i> from the <i><initial value></i> to the <i><final value></i> in turn (using <i><step></i> between each <i><value></i>). The <i><step></i> must be non-zero. If the <i><step></i> is positive, the loop stops when the <i><value></i> becomes larger than the <i><final value></i> . If the <i><step></i> is negative, the loop stops when the <i><value></i> becomes smaller than the <i><final value></i> . The <i><function></i> should absorb one argument.
<hr/> <code>\dim_step_inline:nnnn</code> <hr/> <div>New: 2018-02-18</div>	<code>\dim_step_inline:nnnn {<initial value>} {<step>} {<final value>} {<code>}</code>
	This function first evaluates the <i><initial value></i> , <i><step></i> and <i><final value></i> , all of which should be dimension expressions. Then for each <i><value></i> from the <i><initial value></i> to the <i><final value></i> in turn (using <i><step></i> between each <i><value></i>), the <i><code></i> is inserted into the input stream with #1 replaced by the current <i><value></i> . Thus the <i><code></i> should define a function of one argument (#1).

`\dim_step_variable:nnnNn`
 New: 2018-02-18

`\dim_step_variable:nnnNn`
`{\langle initial value \rangle}{\langle step \rangle}{\langle final value \rangle}{\langle tl var \rangle}{\langle code \rangle}`

This function first evaluates the $\langle initial value \rangle$, $\langle step \rangle$ and $\langle final value \rangle$, all of which should be dimension expressions. Then for each $\langle value \rangle$ from the $\langle initial value \rangle$ to the $\langle final value \rangle$ in turn (using $\langle step \rangle$ between each $\langle value \rangle$), the $\langle code \rangle$ is inserted into the input stream, with the $\langle tl var \rangle$ defined as the current $\langle value \rangle$. Thus the $\langle code \rangle$ should make use of the $\langle tl var \rangle$.

7 Using dim expressions and variables

`\dim_eval:n` ★
 Updated: 2011-10-22

`\dim_eval:n` $\{\langle dimension expression \rangle\}$

Evaluates the $\langle dimension expression \rangle$, expanding any dimensions and token list variables within the $\langle expression \rangle$ to their content (without requiring `\dim_use:N/\tl_use:N`) and applying the standard mathematical rules. The result of the calculation is left in the input stream as a $\langle dimension denotation \rangle$ after two expansions. This is expressed in points (`pt`), and requires suitable termination if used in a TeX-style assignment as it is *not* an $\langle internal dimension \rangle$.

`\dim_sign:n` ★
 New: 2018-11-03

`\dim_sign:n` $\{\langle dimexpr \rangle\}$

Evaluates the $\langle dimexpr \rangle$ then leaves 1 or 0 or -1 in the input stream according to the sign of the result.

`\dim_use:N` ★
`\dim_use:c` ★

`\dim_use:N` $\langle dimension \rangle$

Recovers the content of a $\langle dimension \rangle$ and places it directly in the input stream. An error is raised if the variable does not exist or if it is invalid. Can be omitted in places where a $\langle dimension \rangle$ is required (such as in the argument of `\dim_eval:n`).

TeXhackers note: `\dim_use:N` is the TeX primitive `\the`: this is one of several L^ATeX3 names for this primitive.

`\dim_to_decimal:n` ★
 New: 2014-07-15

`\dim_to_decimal:n` $\{\langle dimexpr \rangle\}$

Evaluates the $\langle dimension expression \rangle$, and leaves the result, expressed in points (`pt`) in the input stream, with *no units*. The result is rounded by TeX to four or five decimal places. If the decimal part of the result is zero, it is omitted, together with the decimal marker.

For example

`\dim_to_decimal:n { 1bp }`

leaves 1.00374 in the input stream, *i.e.* the magnitude of one “big point” when converted to (TeX) points.

<hr/> <code>\dim_to_decimal_in_bp:n</code> ★ <hr/>	<code>\dim_to_decimal_in_bp:n {⟨dimexpr⟩}</code>
<hr/> New: 2014-07-15 <hr/>	Evaluates the <i>⟨dimension expression⟩</i> , and leaves the result, expressed in big points (bp) in the input stream, with <i>no units</i> . The result is rounded by T _E X to four or five decimal places. If the decimal part of the result is zero, it is omitted, together with the decimal marker.

For example

```
\dim_to_decimal_in_bp:n { 1pt }
```

leaves 0.99628 in the input stream, *i.e.* the magnitude of one (T_EX) point when converted to big points.

<hr/> <code>\dim_to_decimal_in_sp:n</code> ★ <hr/>	<code>\dim_to_decimal_in_sp:n {⟨dimexpr⟩}</code>
<hr/> New: 2015-05-18 <hr/>	Evaluates the <i>⟨dimension expression⟩</i> , and leaves the result, expressed in scaled points (sp) in the input stream, with <i>no units</i> . The result is necessarily an integer.

<hr/> <code>\dim_to_decimal_in_unit:nn</code> ★ <hr/>	<code>\dim_to_decimal_in_unit:nn {⟨dimexpr₁⟩} {⟨dimexpr₂⟩}</code>
<hr/> New: 2014-07-15 <hr/>	

Evaluates the *⟨dimension expressions⟩*, and leaves the value of *⟨dimexpr₁⟩*, expressed in a unit given by *⟨dimexpr₂⟩*, in the input stream. The result is a decimal number, rounded by T_EX to four or five decimal places. If the decimal part of the result is zero, it is omitted, together with the decimal marker.

For example

```
\dim_to_decimal_in_unit:nn { 1bp } { 1mm }
```

leaves 0.35277 in the input stream, *i.e.* the magnitude of one big point when converted to millimetres.

Note that this function is not optimised for any particular output and as such may give different results to `\dim_to_decimal_in_bp:n` or `\dim_to_decimal_in_sp:n`. In particular, the latter is able to take a wider range of input values as it is not limited by the ability to calculate a ratio using ε -T_EX primitives, which is required internally by `\dim_to_decimal_in_unit:nn`.

<hr/> <code>\dim_to_fp:n</code> ★ <hr/>	<code>\dim_to_fp:n {⟨dimexpr⟩}</code>
<hr/> New: 2012-05-08 <hr/>	Expands to an internal floating point number equal to the value of the <i>⟨dimexpr⟩</i> in pt. Since dimension expressions are evaluated much faster than their floating point equivalent, <code>\dim_to_fp:n</code> can be used to speed up parts of a computation where a low precision and a smaller range are acceptable.

8 Viewing dim variables

<hr/> <code>\dim_show:N</code> <hr/>	<code>\dim_show:N ⟨dimension⟩</code>
<code>\dim_show:c</code> <hr/>	Displays the value of the <i>⟨dimension⟩</i> on the terminal.

<hr/> <code>\dim_show:n</code> <hr/>	<code>\dim_show:n {⟨<i>dimension expression</i>⟩}</code>
New: 2011-11-22 Updated: 2015-08-07	Displays the result of evaluating the $\langle dimension\ expression \rangle$ on the terminal.
<hr/> <code>\dim_log:N</code> <code>\dim_log:c</code> <hr/>	<code>\dim_log:N ⟨<i>dimension</i>⟩</code>
New: 2014-08-22 Updated: 2015-08-03	Writes the value of the $\langle dimension \rangle$ in the log file.
<hr/> <code>\dim_log:n</code> <hr/>	<code>\dim_log:n {⟨<i>dimension expression</i>⟩}</code>
New: 2014-08-22 Updated: 2015-08-07	Writes the result of evaluating the $\langle dimension\ expression \rangle$ in the log file.

9 Constant dimensions

<hr/> <code>\c_max_dim</code> <hr/>	The maximum value that can be stored as a dimension. This can also be used as a component of a skip.
-------------------------------------	--

<hr/> <code>\c_zero_dim</code> <hr/>	A zero length as a dimension. This can also be used as a component of a skip.
--------------------------------------	---

10 Scratch dimensions

<hr/> <code>\l_tmpa_dim</code> <code>\l_tmppb_dim</code> <hr/>	Scratch dimension for local assignment. These are never used by the kernel code, and so are safe for use with any L ^A T _E X3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.
---	---

<hr/> <code>\g_tmpa_dim</code> <code>\g_tmppb_dim</code> <hr/>	Scratch dimension for global assignment. These are never used by the kernel code, and so are safe for use with any L ^A T _E X3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.
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11 Creating and initialising skip variables

<hr/> <code>\skip_new:N</code> <code>\skip_new:c</code> <hr/>	<code>\skip_new:N ⟨<i>skip</i>⟩</code>
	Creates a new $\langle skip \rangle$ or raises an error if the name is already taken. The declaration is global. The $\langle skip \rangle$ is initially equal to 0 pt.

<code>\skip_const:Nn</code>	<code>\skip_const:Nn <skip> {<skip expression>}</code>
<code>\skip_const:cn</code>	Creates a new constant <i><skip></i> or raises an error if the name is already taken. The value of the <i><skip></i> is set globally to the <i><skip expression></i> .
New: 2012-03-05	

<code>\skip_zero:N</code>	<code>\skip_zero:N <skip></code>
<code>\skip_zero:c</code>	Sets <i><skip></i> to 0 pt.
<code>\skip_gzero:N</code>	
<code>\skip_gzero:c</code>	

<code>\skip_zero_new:N</code>	<code>\skip_zero_new:N <skip></code>
<code>\skip_zero_new:c</code>	Ensures that the <code><skip></code> exists globally by applying <code>\skip_new:N</code> if necessary, then applies <code>\skip_(g)zero:N</code> to leave the <code><skip></code> set to zero.
<code>\skip_gzero_new:N</code>	
<code>\skip_gzero_new:c</code>	
<hr/>	
New: 2012-01-07	

<code>\skip_if_exist_p:N *</code>	<code>\skip_if_exist_p:N $\langle skip \rangle$</code>
<code>\skip_if_exist_p:c *</code>	<code>\skip_if_exist:NTF $\langle skip \rangle$ $\{\langle true\ code \rangle\}$ $\{\langle false\ code \rangle\}$</code>
<code>\skip_if_exist:NTF *</code>	Tests whether the $\langle skip \rangle$ is currently defined. This does not check that the $\langle skip \rangle$ really is a skip variable.
<code>\skip_if_exist:cTF *</code>	

New: 2012-03-03

12 Setting skip variables

<code>\skip_add:Nn</code>	<code>\skip_add:Nn <skip> {<skip expression>}</code>
<code>\skip_add:cn</code>	Adds the result of the <i><skip expression></i> to the current content of the <i><skip></i> .
<code>\skip_gadd:Nn</code>	
<code>\skip_gadd:cn</code>	

Updated: 2011-10-22

<code>\skip_set:Nn</code>	<code>\skip_set:Nn <i><skip></i> {<i><skip expression></i>}</code>
<code>\skip_set:cn</code>	Sets <i><skip></i> to the value of <i><skip expression></i> , which must evaluate to a length with units and may include a rubber component (for example 1 cm plus 0.5 cm).
<code>\skip_gset:Nn</code>	
<code>\skip_gset:cn</code>	
Updated: 2011-10-22	

<code>\skip_set_eq:NN</code>	<code>\skip_set_eq:NN <skip₁₂</code>
<code>\skip_set_eq:(cN Nc cc)</code>	Sets the content of <i><skip_{1 equal to that of <i><skip_{2.}</i>}</i>
<code>\skip_gset_eq:NN</code>	
<code>\skip_gset_eq:(cN Nc cc)</code>	

<code>\skip_sub:Nn</code>	<code>\skip_sub:Nn <skip> {<skip expression>}</code>
<code>\skip_sub:cn</code>	Subtracts the result of the <i><skip expression></i> from the current content of the <i><skip></i> .
<code>\skip_gsub:Nn</code>	
<code>\skip_gsub:cn</code>	

Updated: 2011-10-22

13 Skip expression conditionals

<code>\skip_if_eq_p:nn</code> *	<code>\skip_if_eq_p:nn {\langle skipexpr_1 \rangle} {\langle skipexpr_2 \rangle}</code>
<code>\skip_if_eq:nnTF</code> *	<code>\skip_if_eq:nnTF</code> <code>{\langle skipexpr_1 \rangle} {\langle skipexpr_2 \rangle}</code> <code>{\langle true code \rangle} {\langle false code \rangle}</code>

This function first evaluates each of the $\langle skip\ expressions \rangle$ as described for `\skip_eval:n`. The two results are then compared for exact equality, *i.e.* both the fixed and rubber components must be the same for the test to be true.

<code>\skip_if_finite_p:n</code> *	<code>\skip_if_finite_p:n {\langle skipexpr \rangle}</code>
<code>\skip_if_finite:nTF</code> *	<code>\skip_if_finite:nTF {\langle skipexpr \rangle} {\langle true code \rangle} {\langle false code \rangle}</code>

New: 2012-03-05

Evaluates the $\langle skip\ expression \rangle$ as described for `\skip_eval:n`, and then tests if all of its components are finite.

14 Using skip expressions and variables

<code>\skip_eval:n</code> *	<code>\skip_eval:n {\langle skip expression \rangle}</code>
-----------------------------	---

Updated: 2011-10-22

Evaluates the $\langle skip\ expression \rangle$, expanding any skips and token list variables within the $\langle expression \rangle$ to their content (without requiring `\skip_use:N/\tl_use:N`) and applying the standard mathematical rules. The result of the calculation is left in the input stream as a $\langle glue\ denotation \rangle$ after two expansions. This is expressed in points (`pt`), and requires suitable termination if used in a T_EX-style assignment as it is *not* an $\langle internal\ glue \rangle$.

<code>\skip_use:N</code> *	<code>\skip_use:N \langle skip \rangle</code>
<code>\skip_use:c</code> *	

Recovers the content of a $\langle skip \rangle$ and places it directly in the input stream. An error is raised if the variable does not exist or if it is invalid. Can be omitted in places where a $\langle dimension \rangle$ or $\langle skip \rangle$ is required (such as in the argument of `\skip_eval:n`).

T_EXhackers note: `\skip_use:N` is the T_EX primitive `\the`: this is one of several L^AT_EX3 names for this primitive.

15 Viewing skip variables

<code>\skip_show:N</code>	<code>\skip_show:N \langle skip \rangle</code>
<code>\skip_show:c</code>	

Updated: 2015-08-03

Displays the value of the $\langle skip \rangle$ on the terminal.

<code>\skip_show:n</code>	<code>\skip_show:n {\langle skip expression \rangle}</code>
---------------------------	---

New: 2011-11-22
Updated: 2015-08-07

Displays the result of evaluating the $\langle skip\ expression \rangle$ on the terminal.

<code>\skip_log:N</code>	<code>\skip_log:N <skip></code>
<code>\skip_log:c</code>	Writes the value of the $\langle skip \rangle$ in the log file.
New: 2014-08-22	
Updated: 2015-08-03	

<code>\skip_log:n</code>	<code>\skip_log:n {\langle skip expression \rangle}</code>
	Writes the result of evaluating the $\langle skip expression \rangle$ in the log file.
New: 2014-08-22	
Updated: 2015-08-07	

16 Constant skips

<code>\c_max_skip</code>	The maximum value that can be stored as a skip (equal to <code>\c_max_dim</code> in length), with no stretch nor shrink component.
Updated: 2012-11-02	

<code>\c_zero_skip</code>	A zero length as a skip, with no stretch nor shrink component.
Updated: 2012-11-01	

17 Scratch skips

<code>\l_tmpa_skip</code> <code>\l_tmpb_skip</code>	Scratch skip for local assignment. These are never used by the kernel code, and so are safe for use with any L ^A T _E X3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.
--	--

<code>\g_tmpa_skip</code> <code>\g_tmpb_skip</code>	Scratch skip for global assignment. These are never used by the kernel code, and so are safe for use with any L ^A T _E X3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.
--	---

18 Inserting skips into the output

<code>\skip_horizontal:N</code>	<code>\skip_horizontal:N <skip></code>
<code>\skip_horizontal:c</code>	<code>\skip_horizontal:n {\langle skipexpr \rangle}</code>
<code>\skip_horizontal:n</code>	Inserts a horizontal $\langle skip \rangle$ into the current list. The argument can also be a $\langle dim \rangle$.
Updated: 2011-10-22	
T_EXhackers note: <code>\skip_horizontal:N</code> is the T _E X primitive <code>\hskip</code> renamed.	

<hr/>	
<code>\skip_vertical:N</code>	<code>\skip_vertical:N <skip></code>
<code>\skip_vertical:c</code>	<code>\skip_vertical:n {<skipexpr>}</code>
<code>\skip_vertical:n</code>	Inserts a vertical <code><skip></code> into the current list. The argument can also be a <code><dim></code> .
<hr/> Updated: 2011-10-22 <hr/>	

T_EXhackers note: `\skip_vertical:N` is the T_EX primitive `\vskip` renamed.

19 Creating and initialising muskip variables

<hr/>	
<code>\muskip_new:N</code>	<code>\muskip_new:N <muskip></code>
<code>\muskip_new:c</code>	Creates a new <code><muskip></code> or raises an error if the name is already taken. The declaration is global. The <code><muskip></code> is initially equal to 0 mu.
<hr/>	

<hr/>	
<code>\muskip_const:Nn</code>	<code>\muskip_const:Nn <muskip> {<muskip expression>}</code>
<code>\muskip_const:cn</code>	Creates a new constant <code><muskip></code> or raises an error if the name is already taken. The value of the <code><muskip></code> is set globally to the <code><muskip expression></code> .
<hr/> New: 2012-03-05 <hr/>	

<hr/>	
<code>\muskip_zero:N</code>	<code>\skip_zero:N</code> $\langle muskip \rangle$
<code>\muskip_zero:c</code>	Sets $\langle muskip \rangle$ to 0 mu.
<code>\muskip_gzero:N</code>	
<code>\muskip_gzero:c</code>	
<hr/>	

<code>\muskip_zero_new:N</code>	<code>\muskip_zero_new:N</code> $\langle muskip \rangle$
<code>\muskip_zero_new:c</code>	Ensures that the $\langle muskip \rangle$ exists globally by applying <code>\muskip_new:N</code> if necessary, then applies <code>\muskip_(g)zero:N</code> to leave the $\langle muskip \rangle$ set to zero.
<code>\muskip_gzero_new:N</code>	
<code>\muskip_gzero_new:c</code>	
<hr/>	
New: 2012-01-07	

<hr/>	
<code>\muskip_if_exist_p:N *</code>	<code>\muskip_if_exist_p:N <muskip></code>
<code>\muskip_if_exist_p:c *</code>	<code>\muskip_if_exist:NTF <muskip> {<true code>} {<false code>}</code>
<code>\muskip_if_exist:NTF *</code>	Tests whether the <code><muskip></code> is currently defined. This does not check that the <code><muskip></code> really is a muskip variable.
<code>\muskip_if_exist:cTF *</code>	
<hr/>	
New: 2012-03-03	

20 Setting muskip variables

<code>\muskip_add:Nn</code>	<code>\muskip_add:Nn <muskip> {<muskip expression>}</code>
<code>\muskip_add:cn</code>	Adds the result of the <i><muskip expression></i> to the current content of the <i><muskip></i> .
<code>\muskip_gadd:Nn</code>	
<code>\muskip_gadd:cn</code>	
<hr/>	
Updated: 2011-10-22	

<code>\muskip_set:Nn</code>	<code>\muskip_set:Nn <muskip> {<muskip expression>}</code>
<code>\muskip_set:cn</code>	
<code>\muskip_gset:Nn</code>	Sets <i><muskip></i> to the value of <i><muskip expression></i> , which must evaluate to a math length with units and may include a rubber component (for example 1 mu plus 0.5 mu.
<code>\muskip_gset:cn</code>	
Updated: 2011-10-22	

<code>\muskip_set_eq:NN</code>	<code>\muskip_set_eq:NN <muskip₁> <muskip₂></code>
<code>\muskip_set_eq:(cN Nc cc)</code>	
<code>\muskip_gset_eq:NN</code>	Sets the content of <i><muskip₁></i> equal to that of <i><muskip₂></i> .
<code>\muskip_gset_eq:(cN Nc cc)</code>	

<code>\muskip_sub:Nn</code>	<code>\muskip_sub:Nn <muskip> {<muskip expression>}</code>
<code>\muskip_sub:cn</code>	
<code>\muskip_gsub:Nn</code>	Subtracts the result of the <i><muskip expression></i> from the current content of the <i><muskip></i> .
<code>\muskip_gsub:cn</code>	
Updated: 2011-10-22	

21 Using muskip expressions and variables

<code>\muskip_eval:n *</code>	<code>\muskip_eval:n {<muskip expression>}</code>
Updated: 2011-10-22	Evaluates the <i><muskip expression></i> , expanding any skips and token list variables within the <i><expression></i> to their content (without requiring <code>\muskip_use:N/\tl_use:N</code>) and applying the standard mathematical rules. The result of the calculation is left in the input stream as a <i><mu glue denotation></i> after two expansions. This is expressed in mu , and requires suitable termination if used in a T _E X-style assignment as it is <i>not</i> an <i><internal mu glue></i> .

<code>\muskip_use:N *</code>	<code>\muskip_use:N <muskip></code>
<code>\muskip_use:c *</code>	Recovers the content of a <i><skip></i> and places it directly in the input stream. An error is raised if the variable does not exist or if it is invalid. Can be omitted in places where a <i><dimension></i> is required (such as in the argument of <code>\muskip_eval:n</code>).

T_EXhackers note: `\muskip_use:N` is the T_EX primitive `\the`: this is one of several L^AT_EX3 names for this primitive.

22 Viewing muskip variables

<code>\muskip_show:N</code>	<code>\muskip_show:N <muskip></code>
<code>\muskip_show:c</code>	Displays the value of the <i><muskip></i> on the terminal.
Updated: 2015-08-03	

<hr/> <code>\muskip_show:n</code> <hr/>	<code>\muskip_show:n {⟨<i>muskip expression</i>⟩}</code>
New: 2011-11-22 Updated: 2015-08-07	Displays the result of evaluating the $\langle muskip expression \rangle$ on the terminal.
<hr/> <code>\muskip_log:N</code> <code>\muskip_log:c</code> <hr/>	<code>\muskip_log:N ⟨<i>muskip</i>⟩</code>
New: 2014-08-22 Updated: 2015-08-03	Writes the value of the $\langle muskip \rangle$ in the log file.
<hr/> <code>\muskip_log:n</code> <hr/>	<code>\muskip_log:n {⟨<i>muskip expression</i>⟩}</code>
New: 2014-08-22 Updated: 2015-08-07	Writes the result of evaluating the $\langle muskip expression \rangle$ in the log file.

23 Constant muskips

<hr/> <code>\c_max_muskip</code> <hr/>	The maximum value that can be stored as a muskip, with no stretch nor shrink component.
<hr/> <code>\c_zero_muskip</code> <hr/>	A zero length as a muskip, with no stretch nor shrink component.

24 Scratch muskips

<hr/> <code>\l_tmpa_muskip</code> <code>\l_tmpb_muskip</code> <hr/>	Scratch muskip for local assignment. These are never used by the kernel code, and so are safe for use with any L ^A T _E X3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.
<hr/> <code>\g_tmpa_muskip</code> <code>\g_tmpb_muskip</code> <hr/>	Scratch muskip for global assignment. These are never used by the kernel code, and so are safe for use with any L ^A T _E X3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.

25 Primitive conditional

<hr/> <code>\if_dim:w ★</code> <hr/>	<code>\if_dim:w ⟨<i>dimen</i>₁⟩ ⟨<i>relation</i>⟩ ⟨<i>dimen</i>₂⟩</code> <code> ⟨<i>true code</i>⟩</code> <code> \else:</code> <code> ⟨<i>false</i>⟩</code> <code> \fi:</code>
	Compare two dimensions. The $\langle relation \rangle$ is one of $<$, $=$ or $>$ with category code 12.

T_EXhackers note: This is the T_EX primitive `\ifdim`.

Part XXI

The l3keys package

Key–value interfaces

The key–value method is a popular system for creating large numbers of settings for controlling function or package behaviour. The system normally results in input of the form

```
\MyModuleSetup{
  key-one = value one,
  key-two = value two
}
```

or

```
\MyModuleMacro[
  key-one = value one,
  key-two = value two
]{argument}
```

for the user.

The high level functions here are intended as a method to create key–value controls. Keys are themselves created using a key–value interface, minimising the number of functions and arguments required. Each key is created by setting one or more *properties* of the key:

```
\keys_define:nn { mymodule }
{
  key-one .code:n = code including parameter #1,
  key-two .tl_set:N = \l_mymodule_store_tl
}
```

These values can then be set as with other key–value approaches:

```
\keys_set:nn { mymodule }
{
  key-one = value one,
  key-two = value two
}
```

At a document level, `\keys_set:nn` is used within a document function, for example

```
\DeclareDocumentCommand \MyModuleSetup { m }
{ \keys_set:nn { mymodule } { #1 } }
\DeclareDocumentCommand \MyModuleMacro { o m }
{
  \group_begin:
    \keys_set:nn { mymodule } { #1 }
    % Main code for \MyModuleMacro
  \group_end:
}
```

Key names may contain any tokens, as they are handled internally using `\tl_to_str:n`. As discussed in section 2, it is suggested that the character `/` is reserved for sub-division of keys into logical groups. Functions and variables are *not* expanded when creating key names, and so

```
\tl_set:Nn \l_mymodule_tmp_tl { key }
\keys_define:nn { mymodule }
{
  \l_mymodule_tmp_tl .code:n = code
}
```

creates a key called `\l_mymodule_tmp_tl`, and not one called `key`.

1 Creating keys

`\keys_define:nn`

Updated: 2017-11-14

`\keys_define:nn {<module>} {<keyval list>}`

Parses the *<keyval list>* and defines the keys listed there for *<module>*. The *<module>* name is treated as a string. In practice the *<module>* should be chosen to be unique to the module in question (unless deliberately adding keys to an existing module).

The *<keyval list>* should consist of one or more key names along with an associated key *property*. The properties of a key determine how it acts. The individual properties are described in the following text; a typical use of `\keys_define:nn` might read

```
\keys_define:nn { mymodule }
{
  keyname .code:n = Some~code~using~#1,
  keyname .value_required:n = true
}
```

where the properties of the key begin from the `.` after the key name.

The various properties available take either no arguments at all, or require one or more arguments. This is indicated in the name of the property using an argument specification. In the following discussion, each property is illustrated attached to an arbitrary *<key>*, which when used may be supplied with a *<value>*. All key *definitions* are local.

Key properties are applied in the reading order and so the ordering is significant. Key properties which define “actions”, such as `.code:n`, `.tl_set:N`, *etc.*, override one another. Some other properties are mutually exclusive, notably `.value_required:n` and `.value_forbidden:n`, and so they replace one another. However, properties covering non-exclusive behaviours may be given in any order. Thus for example the following definitions are equivalent.

```
\keys_define:nn { mymodule }
{
  keyname .code:n          = Some~code~using~#1,
  keyname .value_required:n = true
}
\keys_define:nn { mymodule }
{
```

```

    keyname .value_required:n = true,
    keyname .code:n           = Some~code~using~#1
}

```

Note that with the exception of the special `.undefine:` property, all key properties define the key within the current \TeX scope.

```

.bool_set:N
.bool_set:c
.bool_gset:N
.bool_gset:c

```

Updated: 2013-07-08

$\langle key \rangle$ `.bool_set:N = $\langle boolean \rangle$`

Defines $\langle key \rangle$ to set $\langle boolean \rangle$ to $\langle value \rangle$ (which must be either `true` or `false`). If the variable does not exist, it will be created globally at the point that the key is set up.

```

.bool_set_inverse:N
.bool_set_inverse:c
.bool_gset_inverse:N
.bool_gset_inverse:c

```

New: 2011-08-28
Updated: 2013-07-08

$\langle key \rangle$ `.bool_set_inverse:N = $\langle boolean \rangle$`

Defines $\langle key \rangle$ to set $\langle boolean \rangle$ to the logical inverse of $\langle value \rangle$ (which must be either `true` or `false`). If the $\langle boolean \rangle$ does not exist, it will be created globally at the point that the key is set up.

```

.choice:

```

$\langle key \rangle$ `.choice:`

Sets $\langle key \rangle$ to act as a choice key. Each valid choice for $\langle key \rangle$ must then be created, as discussed in section 3.

```

.choices:nn
.choices:(Vn|on|xn)

```

New: 2011-08-21
Updated: 2013-07-10

$\langle key \rangle$ `.choices:nn = { $\langle choices \rangle$ } { $\langle code \rangle$ }`

Sets $\langle key \rangle$ to act as a choice key, and defines a series $\langle choices \rangle$ which are implemented using the $\langle code \rangle$. Inside $\langle code \rangle$, `\l_keys_choice_tl` will be the name of the choice made, and `\l_keys_choice_int` will be the position of the choice in the list of $\langle choices \rangle$ (indexed from 1). Choices are discussed in detail in section 3.

```

.clist_set:N
.clist_set:c
.clist_gset:N
.clist_gset:c

```

New: 2011-09-11

$\langle key \rangle$ `.clist_set:N = $\langle comma list variable \rangle$`

Defines $\langle key \rangle$ to set $\langle comma list variable \rangle$ to $\langle value \rangle$. Spaces around commas and empty items will be stripped. If the variable does not exist, it is created globally at the point that the key is set up.

```

.code:n

```

Updated: 2013-07-10

$\langle key \rangle$ `.code:n = { $\langle code \rangle$ }`

Stores the $\langle code \rangle$ for execution when $\langle key \rangle$ is used. The $\langle code \rangle$ can include one parameter (`#1`), which will be the $\langle value \rangle$ given for the $\langle key \rangle$.

```

.cs_set:Np
.cs_set:cp
.cs_set_protected:Np
.cs_set_protected:cp
.cs_gset:Np
.cs_gset:cp
.cs_gset_protected:Np
.cs_gset_protected:cp

```

New: 2020-01-11

$\langle key \rangle$ `.cs_set:Np = $\langle control sequence \rangle$ $\langle arg. spec. \rangle$`

Defines $\langle key \rangle$ to set $\langle control sequence \rangle$ to have $\langle arg. spec. \rangle$ and replacement text $\langle value \rangle$.

```
.default:n
.default:(V|o|x)
Updated: 2013-07-09
```

$\langle key \rangle$.default:n = { $\langle default \rangle$ }

Creates a $\langle default \rangle$ value for $\langle key \rangle$, which is used if no value is given. This will be used if only the key name is given, but not if a blank $\langle value \rangle$ is given:

```
\keys_define:nn { mymodule }
{
  key .code:n      = Hello~#1,
  key .default:n = World
}
\keys_set:nn { mymodule }
{
  key = Fred, % Prints 'Hello Fred'
  key,      % Prints 'Hello World'
  key = ,    % Prints 'Hello '
}
```

The default does not affect keys where values are required or forbidden. Thus a required value cannot be supplied by a default value, and giving a default value for a key which cannot take a value does not trigger an error.

```
.dim_set:N
.dim_set:c
.dim_gset:N
.dim_gset:c
Updated: 2020-01-17
```

$\langle key \rangle$.dim_set:N = $\langle dimension \rangle$

Defines $\langle key \rangle$ to set $\langle dimension \rangle$ to $\langle value \rangle$ (which must a dimension expression). If the variable does not exist, it is created globally at the point that the key is set up. The key will require a value at point-of-use unless a default is set.

```
.fp_set:N
.fp_set:c
.fp_gset:N
.fp_gset:c
Updated: 2020-01-17
```

$\langle key \rangle$.fp_set:N = $\langle floating point \rangle$

Defines $\langle key \rangle$ to set $\langle floating point \rangle$ to $\langle value \rangle$ (which must a floating point expression). If the variable does not exist, it is created globally at the point that the key is set up. The key will require a value at point-of-use unless a default is set.

```
.groups:n
New: 2013-07-14
```

$\langle key \rangle$.groups:n = { $\langle groups \rangle$ }

Defines $\langle key \rangle$ as belonging to the $\langle groups \rangle$ declared. Groups provide a “secondary axis” for selectively setting keys, and are described in Section 6.

```
.inherit:n
New: 2016-11-22
```

$\langle key \rangle$.inherit:n = { $\langle parents \rangle$ }

Specifies that the $\langle key \rangle$ path should inherit the keys listed as $\langle parents \rangle$. For example, after setting

```
\keys_define:nn { foo } { test .code:n = \tl_show:n {#1} }
\keys_define:nn { } { bar .inherit:n = foo }
```

setting

```
\keys_set:nn { bar } { test = a }
```

will be equivalent to

```
\keys_set:nn { foo } { test = a }
```

<hr/> <code>.initial:n</code> <hr/> <code>.initial:(V o x)</code> <hr/> Updated: 2013-07-09 <hr/>	$\langle key \rangle$ <code>.initial:n = {\langle value \rangle}</code> Initialises the $\langle key \rangle$ with the $\langle value \rangle$, equivalent to $\backslash keys_set:nn \{ \langle module \rangle \} \{ \langle key \rangle = \langle value \rangle \}$
<hr/> <code>.int_set:N</code> <hr/> <code>.int_set:c</code> <hr/> <code>.int_gset:N</code> <hr/> <code>.int_gset:c</code> <hr/> Updated: 2020-01-17 <hr/>	$\langle key \rangle$ <code>.int_set:N = \langle integer \rangle</code> Defines $\langle key \rangle$ to set $\langle integer \rangle$ to $\langle value \rangle$ (which must be an integer expression). If the variable does not exist, it is created globally at the point that the key is set up. The key will require a value at point-of-use unless a default is set.
<hr/> <code>.meta:n</code> <hr/> Updated: 2013-07-10 <hr/>	$\langle key \rangle$ <code>.meta:n = {\langle keyval list \rangle}</code> Makes $\langle key \rangle$ a meta-key, which will set $\langle keyval list \rangle$ in one go. The $\langle keyval list \rangle$ can refer as #1 to the value given at the time the $\langle key \rangle$ is used (or, if no value is given, the $\langle key \rangle$'s default value).
<hr/> <code>.meta:nn</code> <hr/> New: 2013-07-10 <hr/>	$\langle key \rangle$ <code>.meta:nn = {\langle path \rangle} {\langle keyval list \rangle}</code> Makes $\langle key \rangle$ a meta-key, which will set $\langle keyval list \rangle$ in one go using the $\langle path \rangle$ in place of the current one. The $\langle keyval list \rangle$ can refer as #1 to the value given at the time the $\langle key \rangle$ is used (or, if no value is given, the $\langle key \rangle$'s default value).
<hr/> <code>.multichoice:</code> <hr/> New: 2011-08-21 <hr/>	$\langle key \rangle$ <code>.multichoice:</code> Sets $\langle key \rangle$ to act as a multiple choice key. Each valid choice for $\langle key \rangle$ must then be created, as discussed in section 3.
<hr/> <code>.multichoices:nn</code> <hr/> <code>.multichoices:(Vn on xn)</code> <hr/> New: 2011-08-21 Updated: 2013-07-10 <hr/>	$\langle key \rangle$ <code>.multichoices:nn {\langle choices \rangle} {\langle code \rangle}</code> Sets $\langle key \rangle$ to act as a multiple choice key, and defines a series $\langle choices \rangle$ which are implemented using the $\langle code \rangle$. Inside $\langle code \rangle$, <code>\l_keys_choice_tl</code> will be the name of the choice made, and <code>\l_keys_choice_int</code> will be the position of the choice in the list of $\langle choices \rangle$ (indexed from 1). Choices are discussed in detail in section 3.
<hr/> <code>.muskip_set:N</code> <hr/> <code>.muskip_set:c</code> <hr/> <code>.muskip_gset:N</code> <hr/> <code>.muskip_gset:c</code> <hr/> New: 2019-05-05 Updated: 2020-01-17 <hr/>	$\langle key \rangle$ <code>.muskip_set:N = \langle muskip \rangle</code> Defines $\langle key \rangle$ to set $\langle muskip \rangle$ to $\langle value \rangle$ (which must be a muskip expression). If the variable does not exist, it is created globally at the point that the key is set up. The key will require a value at point-of-use unless a default is set.
<hr/> <code>.prop_put:N</code> <hr/> <code>.prop_put:c</code> <hr/> <code>.prop_gput:N</code> <hr/> <code>.prop_gput:c</code> <hr/> New: 2019-01-31 <hr/>	$\langle key \rangle$ <code>.prop_put:N = \langle property list \rangle</code> Defines $\langle key \rangle$ to put the $\langle value \rangle$ onto the $\langle property list \rangle$ stored under the $\langle key \rangle$. If the variable does not exist, it is created globally at the point that the key is set up.

`.skip_set:N` $\langle key \rangle$ `.skip_set:N = $\langle skip \rangle$`

`.skip_set:c`
`.skip_gset:N`
`.skip_gset:c`
Defines $\langle key \rangle$ to set $\langle skip \rangle$ to $\langle value \rangle$ (which must be a skip expression). If the variable does not exist, it is created globally at the point that the key is set up. The key will require a value at point-of-use unless a default is set.

Updated: 2020-01-17

`.tl_set:N` $\langle key \rangle$ `.tl_set:N = $\langle token list variable \rangle$`

`.tl_set:c`
`.tl_gset:N`
`.tl_gset:c`
Defines $\langle key \rangle$ to set $\langle token list variable \rangle$ to $\langle value \rangle$. If the variable does not exist, it is created globally at the point that the key is set up.

`.tl_set_x:N` $\langle key \rangle$ `.tl_set_x:N = $\langle token list variable \rangle$`

`.tl_set_x:c`
`.tl_gset_x:N`
`.tl_gset_x:c`
Defines $\langle key \rangle$ to set $\langle token list variable \rangle$ to $\langle value \rangle$, which will be subjected to an x-type expansion (*i.e.* using `\tl_set:Nx`). If the variable does not exist, it is created globally at the point that the key is set up.

`.undefine:` $\langle key \rangle$ `.undefine:`

New: 2015-07-14

Removes the definition of the $\langle key \rangle$ within the current scope.

`.value_forbidden:n` $\langle key \rangle$ `.value_forbidden:n = true|false`

New: 2015-07-14

Specifies that $\langle key \rangle$ cannot receive a $\langle value \rangle$ when used. If a $\langle value \rangle$ is given then an error will be issued. Setting the property `false` cancels the restriction.

`.value_required:n` $\langle key \rangle$ `.value_required:n = true|false`

New: 2015-07-14

Specifies that $\langle key \rangle$ must receive a $\langle value \rangle$ when used. If a $\langle value \rangle$ is not given then an error will be issued. Setting the property `false` cancels the restriction.

2 Sub-dividing keys

When creating large numbers of keys, it may be desirable to divide them into several sub-groups for a given module. This can be achieved either by adding a sub-division to the module name:

```
\keys_define:nn { mymodule / subgroup }
{ key .code:n = code }
```

or to the key name:

```
\keys_define:nn { mymodule }
{ subgroup / key .code:n = code }
```

As illustrated, the best choice of token for sub-dividing keys in this way is `/`. This is because of the method that is used to represent keys internally. Both of the above code fragments set the same key, which has full name `mymodule/subgroup/key`.

As illustrated in the next section, this subdivision is particularly relevant to making multiple choices.

3 Choice and multiple choice keys

The `l3keys` system supports two types of choice key, in which a series of pre-defined input values are linked to varying implementations. Choice keys are usually created so that the various values are mutually-exclusive: only one can apply at any one time. “Multiple” choice keys are also supported: these allow a selection of values to be chosen at the same time.

Mutually-exclusive choices are created by setting the `.choice:` property:

```
\keys_define:nn { mymodule }
{ key .choice: }
```

For keys which are set up as choices, the valid choices are generated by creating sub-keys of the choice key. This can be carried out in two ways.

In many cases, choices execute similar code which is dependant only on the name of the choice or the position of the choice in the list of all possibilities. Here, the keys can share the same code, and can be rapidly created using the `.choices:nn` property.

```
\keys_define:nn { mymodule }
{
  key .choices:nn =
    { choice-a, choice-b, choice-c }
    {
      You~gave~choice~'\tl_use:N \l_keys_choice_tl',~
      which~is~in~position~\int_use:N \l_keys_choice_int \c_space_tl
      in~the~list.
    }
}
```

The index `\l_keys_choice_int` in the list of choices starts at 1.

`\l_keys_choice_int`
`\l_keys_choice_tl`

Inside the code block for a choice generated using `.choices:nn`, the variables `\l_keys_choice_tl` and `\l_keys_choice_int` are available to indicate the name of the current choice, and its position in the comma list. The position is indexed from 1. Note that, as with standard key code generated using `.code:n`, the value passed to the key (i.e. the choice name) is also available as `#1`.

On the other hand, it is sometimes useful to create choices which use entirely different code from one another. This can be achieved by setting the `.choice:` property of a key, then manually defining sub-keys.

```
\keys_define:nn { mymodule }
{
  key .choice:,
  key / choice-a .code:n = code-a,
  key / choice-b .code:n = code-b,
  key / choice-c .code:n = code-c,
}
```

It is possible to mix the two methods, but manually-created choices should *not* use `\l_keys_choice_tl` or `\l_keys_choice_int`. These variables do not have defined

behaviour when used outside of code created using `.choices:nn` (*i.e.* anything might happen).

It is possible to allow choice keys to take values which have not previously been defined by adding code for the special `unknown` choice. The general behavior of the `unknown` key is described in Section 5. A typical example in the case of a choice would be to issue a custom error message:

```
\keys_define:nn { mymodule }
{
  key .choice:,
  key / choice-a .code:n = code-a,
  key / choice-b .code:n = code-b,
  key / choice-c .code:n = code-c,
  key / unknown .code:n =
    \msg_error:nnxxx { mymodule } { unknown-choice }
    { key } % Name of choice key
    { choice-a , choice-b , choice-c } % Valid choices
    { \exp_not:n {#1} } % Invalid choice given
  %
  %
}
```

Multiple choices are created in a very similar manner to mutually-exclusive choices, using the properties `.multichoice:` and `.multichoices:nn`. As with mutually exclusive choices, multiple choices are define as sub-keys. Thus both

```
\keys_define:nn { mymodule }
{
  key .multichoices:nn =
    { choice-a, choice-b, choice-c }
    {
      You~gave~choice~'\tl_use:N \l_keys_choice_tl',~
      which~is~in~position~
      \int_use:N \l_keys_choice_int \c_space_tl
      in~the~list.
    }
}
```

and

```
\keys_define:nn { mymodule }
{
  key .multichoice:,
  key / choice-a .code:n = code-a,
  key / choice-b .code:n = code-b,
  key / choice-c .code:n = code-c,
}
```

are valid.

When a multiple choice key is set

```
\keys_set:nn { mymodule }
{
  key = { a , b , c } % 'key' defined as a multiple choice
}
```

each choice is applied in turn, equivalent to a `clist` mapping or to applying each value individually:

```
\keys_set:nn { mymodule }
{
  key = a ,
  key = b ,
  key = c ,
}
```

Thus each separate choice will have passed to it the `\l_keys_choice_tl` and `\l_keys_choice_int` in exactly the same way as described for `.choices:nn`.

4 Setting keys

```
\keys_set:nn
\keys_set:(nV|nv|no)
```

Updated: 2017-11-14

```
\keys_set:nn {<module>} {<keyval list>}
```

Parses the `<keyval list>`, and sets those keys which are defined for `<module>`. The behaviour on finding an unknown key can be set by defining a special `unknown` key: this is illustrated later.

```
\l_keys_key_str
\l_keys_path_str
\l_keys_value_tl
```

Updated: 2020-02-08

For each key processed, information of the full *path* of the key, the *name* of the key and the *value* of the key is available within three token list variables. These may be used within the code of the key.

The *value* is everything after the `=`, which may be empty if no value was given. This is stored in `\l_keys_value_tl`, and is not processed in any way by `\keys_set:nn`.

The *path* of the key is a “full” description of the key, and is unique for each key. It consists of the module and full key name, thus for example

```
\keys_set:nn { mymodule } { key-a = some-value }
```

has path `mymodule/key-a` while

```
\keys_set:nn { mymodule } { subset / key-a = some-value }
```

has path `mymodule/subset/key-a`. This information is stored in `\l_keys_path_str`.

The *name* of the key is the part of the path after the last `/`, and thus is not unique. In the preceding examples, both keys have name `key-a` despite having different paths. This information is stored in `\l_keys_key_str`.

5 Handling of unknown keys

If a key has not previously been defined (is unknown), `\keys_set:nn` looks for a special `unknown` key for the same module, and if this is not defined raises an error indicating that

the key name was unknown. This mechanism can be used for example to issue custom error texts.

```
\keys_define:nn { mymodule }
{
  unknown .code:n =
    You~tried~to~set~key~'\l_keys_key_str'~to~'#1'.
}
```

<code>\keys_set_known:nn</code>	<code>\keys_set_known:nn {<module>} {<keyval list>}</code>
<code>\keys_set_known:(nV nv no)</code>	<code>\keys_set_known:nnN {<module>} {<keyval list>} <tl></code>
<code>\keys_set_known:nnN</code>	<code>\keys_set_known:nnnN {<module>} {<keyval list>} {<root>} <tl></code>
<code>\keys_set_known:(nVN nvN noN)</code>	
<code>\keys_set_known:nnnN</code>	
<code>\keys_set_known:(nVnN nvN nonN)</code>	

New: 2011-08-23

Updated: 2019-01-29

These functions set keys which are known for the $\langle module \rangle$, and simply ignore other keys. The `\keys_set_known:nn` function parses the $\langle keyval list \rangle$, and sets those keys which are defined for $\langle module \rangle$. Any keys which are unknown are not processed further by the parser. In addition, `\keys_set_known:nnN` stores the key–value pairs in the $\langle tl \rangle$ in comma-separated form (*i.e.* an edited version of the $\langle keyval list \rangle$). When a $\langle root \rangle$ is given (`\keys_set_known:nnnN`), the key–value entries are returned relative to this point in the key tree. When it is absent, only the key name and value are provided. The correct list is returned by nested calls.

6 Selective key setting

In some cases it may be useful to be able to select only some keys for setting, even though these keys have the same path. For example, with a set of keys defined using

```
\keys_define:nn { mymodule }
{
  key-one .code:n = { \my_func:n {#1} } ,
  key-two .tl_set:N = \l_my_a_tl ,
  key-three .tl_set:N = \l_my_b_tl ,
  key-four .fp_set:N = \l_my_a_fp ,
}
```

the use of `\keys_set:nn` attempts to set all four keys. However, in some contexts it may only be sensible to set some keys, or to control the order of setting. To do this, keys may be assigned to *groups*: arbitrary sets which are independent of the key tree. Thus modifying the example to read

```
\keys_define:nn { mymodule }
{
  key-one .code:n = { \my_func:n {#1} } ,
  key-one .groups:n = { first } ,
  key-two .tl_set:N = \l_my_a_tl ,
}
```

```

key-two    .groups:n = { first }           ,
key-three  .tl_set:N = \l_my_b_tl          ,
key-three  .groups:n = { second }          ,
key-four   .fp_set:N = \l_my_a_fp          ,
}

```

assigns `key-one` and `key-two` to group `first`, `key-three` to group `second`, while `key-four` is not assigned to a group.

Selective key setting may be achieved either by selecting one or more groups to be made “active”, or by marking one or more groups to be ignored in key setting.

<code>\keys_set_filter:nnn</code>	<code>\keys_set_filter:nnn {<module>} {<groups>} {<keyval list>}</code>
<code>\keys_set_filter:(nnV nnv nno)</code>	<code>\keys_set_filter:nnnN {<module>} {<groups>} {<keyval list>} <tl></code>
<code>\keys_set_filter:nnnN</code>	<code>\keys_set_filter:nnnnN {<module>} {<groups>} {<keyval list>} <root></code>
<code>\keys_set_filter:(nnVN nnvN nnoN)</code>	<code><tl></code>
<code>\keys_set_filter:nnnnN</code>	
<code>\keys_set_filter:(nnVnN nnvnN nnonN)</code>	

New: 2013-07-14

Updated: 2019-01-29

Activates key filtering in an “opt-out” sense: keys assigned to any of the `<groups>` specified are ignored. The `<groups>` are given as a comma-separated list. Unknown keys are not assigned to any group and are thus always set. The key–value pairs for each key which is filtered out are stored in the `<tl>` in a comma-separated form (*i.e.* an edited version of the `<keyval list>`). The `\keys_set_filter:nnn` version skips this stage.

Use of `\keys_set_filter:nnnN` can be nested, with the correct residual `<keyval list>` returned at each stage. In the version which takes a `<root>` argument, the key list is returned relative to that point in the key tree. In the cases without a `<root>` argument, only the key names and values are returned.

<code>\keys_set_groups:nnn</code>	<code>\keys_set_groups:nnn {<module>} {<groups>} {<keyval list>}</code>
<code>\keys_set_groups:(nnV nnv nno)</code>	

New: 2013-07-14

Updated: 2017-05-27

Activates key filtering in an “opt-in” sense: only keys assigned to one or more of the `<groups>` specified are set. The `<groups>` are given as a comma-separated list. Unknown keys are not assigned to any group and are thus never set.

7 Utility functions for keys

<code>\keys_if_exist_p:nn *</code>	<code>\keys_if_exist_p:nn {<module>} {<key>}</code>
<code>\keys_if_exist:nnTF *</code>	<code>\keys_if_exist:nnTF {<module>} {<key>} {<true code>} {<false code>}</code>

Updated: 2017-11-14

Tests if the `<key>` exists for `<module>`, *i.e.* if any code has been defined for `<key>`.

<code>\keys_if_choice_exist_p:nnn</code>	<code>*</code>	<code>\keys_if_choice_exist_p:nnn {<module>} {<key>} {<choice>}</code>
<code>\keys_if_choice_exist:nnnTF</code>	<code>*</code>	<code>\keys_if_choice_exist:nnnTF {<module>} {<key>} {<choice>} {<true code>}</code>
		<code>{<false code>}</code>

New: 2011-08-21

Updated: 2017-11-14

Tests if the $\langle choice \rangle$ is defined for the $\langle key \rangle$ within the $\langle module \rangle$, *i.e.* if any code has been defined for $\langle key \rangle / \langle choice \rangle$. The test is **false** if the $\langle key \rangle$ itself is not defined.

<code>\keys_show:nn</code>	<code>\keys_show:nn {<module>} {<key>}</code>
----------------------------	---

Updated: 2015-08-09

Displays in the terminal the information associated to the $\langle key \rangle$ for a $\langle module \rangle$, including the function which is used to actually implement it.

<code>\keys_log:nn</code>	<code>\keys_log:nn {<module>} {<key>}</code>
---------------------------	--

New: 2014-08-22

Updated: 2015-08-09

Writes in the log file the information associated to the $\langle key \rangle$ for a $\langle module \rangle$. See also `\keys_show:nn` which displays the result in the terminal.

8 Low-level interface for parsing key–val lists

To re-cap from earlier, a key–value list is input of the form

```
KeyOne = ValueOne ,
KeyTwo = ValueTwo ,
KeyThree
```

where each key–value pair is separated by a comma from the rest of the list, and each key–value pair does not necessarily contain an equals sign or a value! Processing this type of input correctly requires a number of careful steps, to correctly account for braces, spaces and the category codes of separators.

While the functions described earlier are used as a high-level interface for processing such input, in special circumstances you may wish to use a lower-level approach. The low-level parsing system converts a $\langle key\text{--}value\text{ list} \rangle$ into $\langle keys \rangle$ and associated $\langle values \rangle$. After the parsing phase is completed, the resulting keys and values (or keys alone) are available for further processing. This processing is not carried out by the low-level parser itself, and so the parser requires the names of two functions along with the key–value list. One function is needed to process key–value pairs (it receives two arguments), and a second function is required for keys given without any value (it is called with a single argument).

The parser does not double `#` tokens or expand any input. Active tokens `=` and `,` appearing at the outer level of braces are converted to category “other” (12) so that the parser does not “miss” any due to category code changes. Spaces are removed from the ends of the keys and values. Keys and values which are given in braces have exactly one set removed (after space trimming), thus

```
key = {value here},
```

and

```
key = value here,
```

are treated identically.

`\keyval_parse:nnn` ☆

New: 2020-12-19

`\keyval_parse:nnn {⟨code1⟩} {⟨code2⟩} {⟨key-value list⟩}`

Parses the *⟨key-value list⟩* into a series of *⟨keys⟩* and associated *⟨values⟩*, or keys alone (if no *⟨value⟩* was given). *⟨code₁⟩* receives each *⟨key⟩* (with no *⟨value⟩*) as a trailing brace group, whereas *⟨code₂⟩* is appended by two brace groups, the *⟨key⟩* and *⟨value⟩*. The order of the *⟨keys⟩* in the *⟨key-value list⟩* is preserved. Thus

```
\keyval_parse:nnn
{ \use_none:nn { code 1 } }
{ \use_none:nnn { code 2 } }
{ key1 = value1 , key2 = value2, key3 = , key4 }
```

is converted into an input stream

```
\use_none:nnn { code 2 } { key1 } { value1 }
\use_none:nnn { code 2 } { key2 } { value2 }
\use_none:nnn { code 2 } { key3 } { }
\use_none:nn { code 1 } { key4 }
```

Note that there is a difference between an empty value (an equals sign followed by nothing) and a missing value (no equals sign at all). Spaces are trimmed from the ends of the *⟨key⟩* and *⟨value⟩*, then one *outer* set of braces is removed from the *⟨key⟩* and *⟨value⟩* as part of the processing. If you need exactly the output shown above, you'll need to either **x**-type or **e**-type expand the function.

TeXhackers note: The result of each list element is returned within `\exp_not:n`, which means that the converted input stream does not expand further when appearing in an **x**-type or **e**-type argument expansion.

`\keyval_parse:NNn` ☆

Updated: 2020-12-19

`\keyval_parse:NNn` $\langle function_1 \rangle$ $\langle function_2 \rangle$ $\{ \langle key-value list \rangle \}$

Parses the $\langle key-value list \rangle$ into a series of $\langle keys \rangle$ and associated $\langle values \rangle$, or keys alone (if no $\langle value \rangle$ was given). $\langle function_1 \rangle$ should take one argument, while $\langle function_2 \rangle$ should absorb two arguments. After `\keyval_parse:NNn` has parsed the $\langle key-value list \rangle$, $\langle function_1 \rangle$ is used to process keys given with no value and $\langle function_2 \rangle$ is used to process keys given with a value. The order of the $\langle keys \rangle$ in the $\langle key-value list \rangle$ is preserved. Thus

```
\keyval_parse:NNn \function:n \function:nn
{ key1 = value1 , key2 = value2, key3 = , key4 }
```

is converted into an input stream

```
\function:nn { key1 } { value1 }
\function:nn { key2 } { value2 }
\function:nn { key3 } { }
\function:n { key4 }
```

Note that there is a difference between an empty value (an equals sign followed by nothing) and a missing value (no equals sign at all). Spaces are trimmed from the ends of the $\langle key \rangle$ and $\langle value \rangle$, then one *outer* set of braces is removed from the $\langle key \rangle$ and $\langle value \rangle$ as part of the processing.

This shares the implementation of `\keyval_parse:nnn`, the difference is only semantically.

TeXhackers note: The result is returned within `\exp_not:n`, which means that the converted input stream does not expand further when appearing in an **x**-type or **e**-type argument expansion.

Part XXII

The l3intarray package: fast global integer arrays

1 l3intarray documentation

For applications requiring heavy use of integers, this module provides arrays which can be accessed in constant time (contrast l3seq, where access time is linear). These arrays have several important features

- The size of the array is fixed and must be given at point of initialisation
- The absolute value of each entry has maximum $2^{30} - 1$ (*i.e.* one power lower than the usual `\c_max_int` ceiling of $2^{31} - 1$)

The use of `intarray` data is therefore recommended for cases where the need for fast access is of paramount importance.

`\intarray_new:Nn`
`\intarray_new:cn`

New: 2018-03-29

`\intarray_new:Nn` $\langle\textit{intarray var}\rangle$ $\{\langle\textit{size}\rangle\}$

Evaluates the integer expression $\langle\textit{size}\rangle$ and allocates an $\langle\textit{integer array variable}\rangle$ with that number of (zero) entries. The variable name should start with `\g_` because assignments are always global.

`\intarray_count:N *`
`\intarray_count:c *`

New: 2018-03-29

`\intarray_count:N` $\langle\textit{intarray var}\rangle$

Expands to the number of entries in the $\langle\textit{integer array variable}\rangle$. Contrarily to `\seq-count:N` this is performed in constant time.

`\intarray_gset:Nnn`
`\intarray_gset:cnn`

New: 2018-03-29

`\intarray_gset:Nnn` $\langle\textit{intarray var}\rangle$ $\{\langle\textit{position}\rangle\}$ $\{\langle\textit{value}\rangle\}$

Stores the result of evaluating the integer expression $\langle\textit{value}\rangle$ into the $\langle\textit{integer array variable}\rangle$ at the (integer expression) $\langle\textit{position}\rangle$. If the $\langle\textit{position}\rangle$ is not between 1 and the `\intarray_count:N`, or the $\langle\textit{value}\rangle$'s absolute value is bigger than $2^{30} - 1$, an error occurs. Assignments are always global.

`\intarray_const_from_clist:Nn`
`\intarray_const_from_clist:cn`

New: 2018-05-04

`\intarray_const_from_clist:Nn` $\langle\textit{intarray var}\rangle$ $\langle\textit{intexpr clist}\rangle$

Creates a new constant $\langle\textit{integer array variable}\rangle$ or raises an error if the name is already taken. The $\langle\textit{integer array variable}\rangle$ is set (globally) to contain as its items the results of evaluating each $\langle\textit{integer expression}\rangle$ in the $\langle\textit{comma list}\rangle$.

`\intarray_gzero:N`
`\intarray_gzero:c`

New: 2018-05-04

`\intarray_gzero:N` $\langle\textit{intarray var}\rangle$

Sets all entries of the $\langle\textit{integer array variable}\rangle$ to zero. Assignments are always global.

<hr/>	
<code>\intarray_item:Nn</code> *	<code>\intarray_item:Nn <intarray var> {<position>}</code>
<code>\intarray_item:cn</code> *	Expands to the integer entry stored at the (integer expression) <i><position></i> in the <i><integer array variable></i> . If the <i><position></i> is not between 1 and the <code>\intarray_count:N</code> , an error occurs.
<hr/>	
New: 2018-03-29	
<hr/>	
<code>\intarray_rand_item:N</code> *	<code>\intarray_rand_item:N <intarray var></code>
<code>\intarray_rand_item:c</code> *	Selects a pseudo-random item of the <i><integer array></i> . If the <i><integer array></i> is empty, produce an error.
<hr/>	
New: 2018-05-05	
<hr/>	
<code>\intarray_show:N</code>	<code>\intarray_show:N <intarray var></code>
<code>\intarray_show:c</code>	<code>\intarray_log:N <intarray var></code>
<code>\intarray_log:N</code>	Displays the items in the <i><integer array variable></i> in the terminal or writes them in the log file.
<code>\intarray_log:c</code>	
<hr/>	
New: 2018-05-04	
<hr/>	

1.1 Implementation notes

It is a wrapper around the `\fontdimen` primitive, used to store arrays of integers (with a restricted range: absolute value at most $2^{30} - 1$). In contrast to `l3seq` sequences the access to individual entries is done in constant time rather than linear time, but only integers can be stored. More precisely, the primitive `\fontdimen` stores dimensions but the `l3intarray` package transparently converts these from/to integers. Assignments are always global.

While LuaTeX’s memory is extensible, other engines can “only” deal with a bit less than 4×10^6 entries in all `\fontdimen` arrays combined (with default TeX Live settings).

Part XXIII

The l3fp package: Floating points

A decimal floating point number is one which is stored as a significand and a separate exponent. The module implements expandably a wide set of arithmetic, trigonometric, and other operations on decimal floating point numbers, to be used within floating point expressions. Floating point expressions support the following operations with their usual precedence.

- Basic arithmetic: addition $x + y$, subtraction $x - y$, multiplication $x * y$, division x / y , square root \sqrt{x} , and parentheses.
 - Comparison operators: $x < y$, $x \leq y$, $x > y$, $x != y$ etc.
 - Boolean logic: sign $\text{sign } x$, negation $!x$, conjunction $x \&\& y$, disjunction $x || y$, ternary operator $x ? y : z$.
 - Exponentials: $\exp x$, $\ln x$, x^y , $\log b x$.
 - Integer factorial: $\text{fact } x$.
 - Trigonometry: $\sin x$, $\cos x$, $\tan x$, $\cot x$, $\sec x$, $\csc x$ expecting their arguments in radians, and $\text{sind } x$, $\text{cosd } x$, $\text{tand } x$, $\text{cotd } x$, $\text{secd } x$, $\text{cscd } x$ expecting their arguments in degrees.
 - Inverse trigonometric functions: $\text{asin } x$, $\text{acos } x$, $\text{atan } x$, $\text{acot } x$, $\text{asec } x$, $\text{acsc } x$ giving a result in radians, and $\text{asind } x$, $\text{acosd } x$, $\text{atand } x$, $\text{acotd } x$, $\text{asecd } x$, $\text{acscd } x$ giving a result in degrees.
- (not yet) Hyperbolic functions and their inverse functions: $\sinh x$, $\cosh x$, $\tanh x$, $\coth x$, $\text{sech } x$, $\text{csch } x$, and $\text{asinh } x$, $\text{acosh } x$, $\text{atanh } x$, $\text{acoth } x$, $\text{asech } x$, $\text{acsch } x$.
- Extrema: $\max(x_1, x_2, \dots)$, $\min(x_1, x_2, \dots)$, $\text{abs}(x)$.
 - Rounding functions, controlled by two optional values, n (number of places, 0 by default) and t (behavior on a tie, NaN by default):
 - $\text{trunc}(x, n)$ rounds towards zero,
 - $\text{floor}(x, n)$ rounds towards $-\infty$,
 - $\text{ceil}(x, n)$ rounds towards $+\infty$,
 - $\text{round}(x, n, t)$ rounds to the closest value, with ties rounded to an even value by default, towards zero if $t = 0$, towards $+\infty$ if $t > 0$ and towards $-\infty$ if $t < 0$.

And (not yet) modulo, and “quantize”.

- Random numbers: $\text{rand}()$, $\text{randint}(m, n)$.
- Constants: pi , deg (one degree in radians).
- Dimensions, automatically expressed in points, e.g., pc is 12.

- Automatic conversion (no need for `\langle type \rangle_use:N`) of integer, dimension, and skip variables to floating point numbers, expressing dimensions in points and ignoring the stretch and shrink components of skips.
- Tuples: (x_1, \dots, x_n) that can be stored in variables, added together, multiplied or divided by a floating point number, and nested.

Floating point numbers can be given either explicitly (in a form such as $1.234\text{e-}34$, or $-.0001$), or as a stored floating point variable, which is automatically replaced by its current value. A “floating point” is a floating point number or a tuple thereof. See section 9.1 for a description of what a floating point is, section 9.2 for details about how an expression is parsed, and section 9.3 to know what the various operations do. Some operations may raise exceptions (error messages), described in section 7.

An example of use could be the following.

```
\LaTeX{} can now compute: $ \frac{\sin (3.5)}{2} + 2\cdot 10^{-3}
= \ExplSyntaxOn \fp_to_decimal:n {\sin(3.5)/2 + 2e-3} $.
```

The operation `round` can be used to limit the result’s precision. Adding `+0` avoids the possibly undesirable output `-0`, replacing it by `+0`. However, the `l3fp` module is mostly meant as an underlying tool for higher-level commands. For example, one could provide a function to typeset nicely the result of floating point computations.

```
\documentclass{article}
\usepackage{xparse, siunitx}
\ExplSyntaxOn
\NewDocumentCommand { \calcnun } { m }
{ \num { \fp_to_scientific:n {#1} } }
\ExplSyntaxOff
\begin{document}
\calcnun { 2 pi * sin ( 2.3 ^ 5 ) }
\end{document}
```

See the documentation of `siunitx` for various options of `\num`.

1 Creating and initialising floating point variables

<code>\fp_new:N</code>	<code>\fp_new:N <fp var></code>
<code>\fp_new:c</code>	Creates a new <code><fp var></code> or raises an error if the name is already taken. The declaration is global. The <code><fp var></code> is initially <code>+0</code> .
Updated: 2012-05-08	
<code>\fp_const:Nn</code>	<code>\fp_const:Nn <fp var> {<floating point expression>}</code>
<code>\fp_const:cn</code>	Creates a new constant <code><fp var></code> or raises an error if the name is already taken. The <code><fp var></code> is set globally equal to the result of evaluating the <code><floating point expression></code> .
Updated: 2012-05-08	
<code>\fp_zero:N</code>	<code>\fp_zero:N <fp var></code>
<code>\fp_zero:c</code>	Sets the <code><fp var></code> to <code>+0</code> .
<code>\fp_gzero:N</code>	
<code>\fp_gzero:c</code>	
Updated: 2012-05-08	

```
\fp_zero_new:N
\fp_zero_new:c
\fp_gzero_new:N
\fp_gzero_new:c
```

Updated: 2012-05-08

```
\fp_zero_new:N <fp var>
```

Ensures that the $\langle fp\ var \rangle$ exists globally by applying $\backslash fp_new:N$ if necessary, then applies $\backslash fp_(\mathit{g})zero:N$ to leave the $\langle fp\ var \rangle$ set to +0.

2 Setting floating point variables

```
\fp_set:Nn
\fp_set:cn
\fp_gset:Nn
\fp_gset:cn
```

Updated: 2012-05-08

```
\fp_set:Nn <fp var> {<floating point expression>}
```

Sets $\langle fp\ var \rangle$ equal to the result of computing the $\langle floating\ point\ expression \rangle$.

```
\fp_set_eq:NN
\fp_set_eq:(cN|Nc|cc)
\fp_gset_eq:NN
\fp_gset_eq:(cN|Nc|cc)
```

Updated: 2012-05-08

```
\fp_set_eq:NN <fp var1> <fp var2>
```

Sets the floating point variable $\langle fp\ var_1 \rangle$ equal to the current value of $\langle fp\ var_2 \rangle$.

```
\fp_add:Nn
\fp_add:cn
\fp_gadd:Nn
\fp_gadd:cn
```

Updated: 2012-05-08

```
\fp_add:Nn <fp var> {<floating point expression>}
```

Adds the result of computing the $\langle floating\ point\ expression \rangle$ to the $\langle fp\ var \rangle$. This also applies if $\langle fp\ var \rangle$ and $\langle floating\ point\ expression \rangle$ evaluate to tuples of the same size.

```
\fp_sub:Nn
\fp_sub:cn
\fp_gsub:Nn
\fp_gsub:cn
```

Updated: 2012-05-08

```
\fp_sub:Nn <fp var> {<floating point expression>}
```

Subtracts the result of computing the $\langle floating\ point\ expression \rangle$ from the $\langle fp\ var \rangle$. This also applies if $\langle fp\ var \rangle$ and $\langle floating\ point\ expression \rangle$ evaluate to tuples of the same size.

3 Using floating points

```
\fp_eval:n ★
```

New: 2012-05-08
Updated: 2012-07-08

```
\fp_eval:n {<floating point expression>}
```

Evaluates the $\langle floating\ point\ expression \rangle$ and expresses the result as a decimal number with no exponent. Leading or trailing zeros may be inserted to compensate for the exponent. Non-significant trailing zeros are trimmed, and integers are expressed without a decimal separator. The values $\pm\infty$ and NaN trigger an “invalid operation” exception. For a tuple, each item is converted using $\backslash fp_eval:n$ and they are combined as $(\langle fp_1 \rangle, \sqcup \langle fp_2 \rangle, \sqcup \dots \langle fp_n \rangle)$ if $n > 1$ and $(\langle fp_1 \rangle,)$ or $()$ for fewer items. This function is identical to $\backslash fp_to_decimal:n$.

<code>\fp_sign:n</code> *	<code>\fp_sign:n {<fpexpr>}</code>
---------------------------	--

New: 2018-11-03

Evaluates the $\langle fpexpr \rangle$ and leaves its sign in the input stream using `\fp_eval:n {sign(<result>)}`: +1 for positive numbers and for $+\infty$, -1 for negative numbers and for $-\infty$, ± 0 for ± 0 . If the operand is a tuple or is NaN, then “invalid operation” occurs and the result is 0.

<code>\fp_to_decimal:N</code> *	<code>\fp_to_decimal:N <fp var></code>
<code>\fp_to_decimal:c</code> *	<code>\fp_to_decimal:n {<floating point expression>}</code>
<code>\fp_to_decimal:n</code> *	

New: 2012-05-08

Updated: 2012-07-08

Evaluates the $\langle floating point expression \rangle$ and expresses the result as a decimal number with no exponent. Leading or trailing zeros may be inserted to compensate for the exponent. Non-significant trailing zeros are trimmed, and integers are expressed without a decimal separator. The values $\pm\infty$ and NaN trigger an “invalid operation” exception. For a tuple, each item is converted using `\fp_to_decimal:n` and they are combined as $\langle fp_1 \rangle, \sqcup \langle fp_2 \rangle, \sqcup \dots \langle fp_n \rangle$ if $n > 1$ and $\langle fp_1 \rangle,)$ or $()$ for fewer items.

<code>\fp_to_dim:N</code> *	<code>\fp_to_dim:N <fp var></code>
<code>\fp_to_dim:c</code> *	<code>\fp_to_dim:n {<floating point expression>}</code>
<code>\fp_to_dim:n</code> *	

Updated: 2016-03-22

Evaluates the $\langle floating point expression \rangle$ and expresses the result as a dimension (in pt) suitable for use in dimension expressions. The output is identical to `\fp_to_decimal:n`, with an additional trailing pt (both letter tokens). In particular, the result may be outside the range $[-2^{14} + 2^{-17}, 2^{14} - 2^{-17}]$ of valid T_EX dimensions, leading to overflow errors if used as a dimension. Tuples, as well as the values $\pm\infty$ and NaN, trigger an “invalid operation” exception.

<code>\fp_to_int:N</code> *	<code>\fp_to_int:N <fp var></code>
<code>\fp_to_int:c</code> *	<code>\fp_to_int:n {<floating point expression>}</code>
<code>\fp_to_int:n</code> *	

Updated: 2012-07-08

Evaluates the $\langle floating point expression \rangle$, and rounds the result to the closest integer, rounding exact ties to an even integer. The result may be outside the range $[-2^{31} + 1, 2^{31} - 1]$ of valid T_EX integers, leading to overflow errors if used in an integer expression. Tuples, as well as the values $\pm\infty$ and NaN, trigger an “invalid operation” exception.

<code>\fp_to_scientific:N</code> *	<code>\fp_to_scientific:N <fp var></code>
<code>\fp_to_scientific:c</code> *	<code>\fp_to_scientific:n {<floating point expression>}</code>
<code>\fp_to_scientific:n</code> *	

New: 2012-05-08

Updated: 2016-03-22

Evaluates the $\langle floating point expression \rangle$ and expresses the result in scientific notation:

$\langle optional - \rangle \langle digit \rangle . \langle 15 digits \rangle e \langle optional sign \rangle \langle exponent \rangle$

The leading $\langle digit \rangle$ is non-zero except in the case of ± 0 . The values $\pm\infty$ and NaN trigger an “invalid operation” exception. Normal category codes apply: thus the `e` is category code 11 (a letter). For a tuple, each item is converted using `\fp_to_scientific:n` and they are combined as $\langle fp_1 \rangle, \sqcup \langle fp_2 \rangle, \sqcup \dots \langle fp_n \rangle$ if $n > 1$ and $\langle fp_1 \rangle,)$ or $()$ for fewer items.

<code>\fp_to_tl:N</code>	★	<code>\fp_to_tl:N <fp var></code>
<code>\fp_to_tl:c</code>	★	<code>\fp_to_tl:n {\floating point expression}</code>
<code>\fp_to_tl:n</code>	★	Evaluates the <i><floating point expression></i> and expresses the result in (almost) the shortest possible form. Numbers in the ranges $(0, 10^{-3})$ and $[10^{16}, \infty)$ are expressed in scientific notation with trailing zeros trimmed and no decimal separator when there is a single significant digit (this differs from <code>\fp_to_scientific:n</code>). Numbers in the range $[10^{-3}, 10^{16})$ are expressed in a decimal notation without exponent, with trailing zeros trimmed, and no decimal separator for integer values (see <code>\fp_to_decimal:n</code> . Negative numbers start with <code>-</code> . The special values ± 0 , $\pm \infty$ and NaN are rendered as <code>0</code> , <code>-0</code> , <code>inf</code> , <code>-inf</code> , and <code>nan</code> respectively. Normal category codes apply and thus <code>inf</code> or <code>nan</code> , if produced, are made up of letters. For a tuple, each item is converted using <code>\fp_to_tl:n</code> and they are combined as $\langle \langle fp_1 \rangle, \sqcup \langle fp_2 \rangle, \sqcup \dots \langle fp_n \rangle \rangle$ if $n > 1$ and $\langle \langle fp_1 \rangle, \rangle$ or $\langle \rangle$ for fewer items.

Updated: 2016-03-22

<code>\fp_use:N</code>	★	<code>\fp_use:N <fp var></code>
<code>\fp_use:c</code>	★	Inserts the value of the <i><fp var></i> into the input stream as a decimal number with no exponent. Leading or trailing zeros may be inserted to compensate for the exponent. Non-significant trailing zeros are trimmed. Integers are expressed without a decimal separator. The values $\pm \infty$ and NaN trigger an “invalid operation” exception. For a tuple, each item is converted using <code>\fp_to_decimal:n</code> and they are combined as $\langle \langle fp_1 \rangle, \sqcup \langle fp_2 \rangle, \sqcup \dots \langle fp_n \rangle \rangle$ if $n > 1$ and $\langle \langle fp_1 \rangle, \rangle$ or $\langle \rangle$ for fewer items. This function is identical to <code>\fp_to_decimal:N</code> .

Updated: 2012-07-08

4 Floating point conditionals

<code>\fp_if_exist_p:N</code>	★	<code>\fp_if_exist_p:N <fp var></code>
<code>\fp_if_exist_p:c</code>	★	<code>\fp_if_exist:N\TF <fp var> {\true code} {\false code}</code>
<code>\fp_if_exist:N\TF</code>	★	Tests whether the <i><fp var></i> is currently defined. This does not check that the <i><fp var></i> really is a floating point variable.
<code>\fp_if_exist:c\TF</code>	★	

Updated: 2012-05-08

<code>\fp_compare_p:nNn</code>	★	<code>\fp_compare_p:nNn {\fpexpr₁} <relation> {\fpexpr₂}</code>
<code>\fp_compare:nNn\TF</code>	★	<code>\fp_compare:nNn\TF {\fpexpr₁} <relation> {\fpexpr₂} {\true code} {\false code}</code>

Updated: 2012-05-08

Compares the *<fpexpr₁>* and the *<fpexpr₂>*, and returns `true` if the *<relation>* is obeyed. Two floating points x and y may obey four mutually exclusive relations: $x < y$, $x = y$, $x > y$, or $x?y$ (“not ordered”). The last case occurs exactly if one or both operands is NaN or is a tuple, unless they are equal tuples. Note that a NaN is distinct from any value, even another NaN, hence $x = x$ is not true for a NaN. To test if a value is NaN, compare it to an arbitrary number with the “not ordered” relation.

```

\fp_compare:nNn\TF { <value> } ? { 0 }
{ } % <value> is nan
{ } % <value> is not nan

```

Tuples are equal if they have the same number of items and items compare equal (in particular there must be no NaN). At present any other comparison with tuples yields ? (not ordered). This is experimental.

This function is less flexible than `\fp_compare:n\TF` but slightly faster. It is provided for consistency with `\int_compare:nNn\TF` and `\dim_compare:nNn\TF`.

<code>\fp_compare_p:n</code> ☆	<code>\fp_compare_p:n</code>
<code>\fp_compare:nTF</code> ☆	{
Updated: 2013-12-14	$\langle fpexpr_1 \rangle$ $\langle relation_1 \rangle$
	...
	$\langle fpexpr_N \rangle$ $\langle relation_N \rangle$
	$\langle fpexpr_{N+1} \rangle$
	}
	<code>\fp_compare:nTF</code>
	{
	$\langle fpexpr_1 \rangle$ $\langle relation_1 \rangle$
	...
	$\langle fpexpr_N \rangle$ $\langle relation_N \rangle$
	$\langle fpexpr_{N+1} \rangle$
	}
	{ $\langle true\ code \rangle$ } { $\langle false\ code \rangle$ }

Evaluates the $\langle floating\ point\ expressions \rangle$ as described for `\fp_eval:n` and compares consecutive result using the corresponding $\langle relation \rangle$, namely it compares $\langle fpexpr_1 \rangle$ and $\langle fpexpr_2 \rangle$ using the $\langle relation_1 \rangle$, then $\langle fpexpr_2 \rangle$ and $\langle fpexpr_3 \rangle$ using the $\langle relation_2 \rangle$, until finally comparing $\langle fpexpr_N \rangle$ and $\langle fpexpr_{N+1} \rangle$ using the $\langle relation_N \rangle$. The test yields **true** if all comparisons are **true**. Each $\langle floating\ point\ expression \rangle$ is evaluated only once. Contrarily to `\int_compare:nTF`, all $\langle floating\ point\ expressions \rangle$ are computed, even if one comparison is **false**. Two floating points x and y may obey four mutually exclusive relations: $x < y$, $x = y$, $x > y$, or $x?y$ (“not ordered”). The last case occurs exactly if one or both operands is NaN or is a tuple, unless they are equal tuples. Each $\langle relation \rangle$ can be any (non-empty) combination of $<$, $=$, $>$, and $?$, plus an optional leading $!$ (which negates the $\langle relation \rangle$), with the restriction that the $\langle relation \rangle$ may not start with $?$, as this symbol has a different meaning (in combination with $:$) within floating point expressions. The comparison $x \langle relation \rangle y$ is then **true** if the $\langle relation \rangle$ does not start with $!$ and the actual relation ($<$, $=$, $>$, or $?$) between x and y appears within the $\langle relation \rangle$, or on the contrary if the $\langle relation \rangle$ starts with $!$ and the relation between x and y does not appear within the $\langle relation \rangle$. Common choices of $\langle relation \rangle$ include $>=$ (greater or equal), $!=$ (not equal), $!?$ or $<=>$ (comparable).

This function is more flexible than `\fp_compare:nNnTF` and only slightly slower.

5 Floating point expression loops

<code>\fp_do_until:nNnn</code> ☆	<code>\fp_do_until:nNnn {$\langle fpexpr_1 \rangle$} $\langle relation \rangle$ {$\langle fpexpr_2 \rangle$} {$\langle code \rangle$}</code>
New: 2012-08-16	Places the $\langle code \rangle$ in the input stream for \TeX to process, and then evaluates the relationship between the two $\langle floating\ point\ expressions \rangle$ as described for <code>\fp_compare:nNnTF</code> . If the test is false then the $\langle code \rangle$ is inserted into the input stream again and a loop occurs until the $\langle relation \rangle$ is true .
<code>\fp_do_while:nNnn</code> ☆	<code>\fp_do_while:nNnn {$\langle fpexpr_1 \rangle$} $\langle relation \rangle$ {$\langle fpexpr_2 \rangle$} {$\langle code \rangle$}</code>
New: 2012-08-16	Places the $\langle code \rangle$ in the input stream for \TeX to process, and then evaluates the relationship between the two $\langle floating\ point\ expressions \rangle$ as described for <code>\fp_compare:nNnTF</code> . If the test is true then the $\langle code \rangle$ is inserted into the input stream again and a loop occurs until the $\langle relation \rangle$ is false .

<hr/>	
<code>\fp_until_do:nNnn</code> ☆	<code>\fp_until_do:nNnn {<fpexpr1>} <relation> {<fpexpr2>} {<code>}</code>
New: 2012-08-16	Evaluates the relationship between the two <i><floating point expressions></i> as described for <code>\fp_compare:nNnTF</code> , and then places the <i><code></i> in the input stream if the <i><relation></i> is false. After the <i><code></i> has been processed by T _E X the test is repeated, and a loop occurs until the test is true .
<hr/>	
<code>\fp_while_do:nNnn</code> ☆	<code>\fp_while_do:nNnn {<fpexpr1>} <relation> {<fpexpr2>} {<code>}</code>
New: 2012-08-16	Evaluates the relationship between the two <i><floating point expressions></i> as described for <code>\fp_compare:nNnTF</code> , and then places the <i><code></i> in the input stream if the <i><relation></i> is true. After the <i><code></i> has been processed by T _E X the test is repeated, and a loop occurs until the test is false .
<hr/>	
<code>\fp_do_until:nn</code> ☆	<code>\fp_do_until:nn { <fpexpr1> <relation> <fpexpr2> } {<code>}</code>
New: 2012-08-16 Updated: 2013-12-14	Places the <i><code></i> in the input stream for T _E X to process, and then evaluates the relationship between the two <i><floating point expressions></i> as described for <code>\fp_compare:nTF</code> . If the test is false then the <i><code></i> is inserted into the input stream again and a loop occurs until the <i><relation></i> is true .
<hr/>	
<code>\fp_do_while:nn</code> ☆	<code>\fp_do_while:nn { <fpexpr1> <relation> <fpexpr2> } {<code>}</code>
New: 2012-08-16 Updated: 2013-12-14	Places the <i><code></i> in the input stream for T _E X to process, and then evaluates the relationship between the two <i><floating point expressions></i> as described for <code>\fp_compare:nTF</code> . If the test is true then the <i><code></i> is inserted into the input stream again and a loop occurs until the <i><relation></i> is false .
<hr/>	
<code>\fp_until_do:nn</code> ☆	<code>\fp_until_do:nn { <fpexpr1> <relation> <fpexpr2> } {<code>}</code>
New: 2012-08-16 Updated: 2013-12-14	Evaluates the relationship between the two <i><floating point expressions></i> as described for <code>\fp_compare:nTF</code> , and then places the <i><code></i> in the input stream if the <i><relation></i> is false. After the <i><code></i> has been processed by T _E X the test is repeated, and a loop occurs until the test is true .
<hr/>	
<code>\fp_while_do:nn</code> ☆	<code>\fp_while_do:nn { <fpexpr1> <relation> <fpexpr2> } {<code>}</code>
New: 2012-08-16 Updated: 2013-12-14	Evaluates the relationship between the two <i><floating point expressions></i> as described for <code>\fp_compare:nTF</code> , and then places the <i><code></i> in the input stream if the <i><relation></i> is true. After the <i><code></i> has been processed by T _E X the test is repeated, and a loop occurs until the test is false .

`\fp_step_function:nnnN` ☆
`\fp_step_function:nnnc` ☆

New: 2016-11-21
Updated: 2016-12-06

`\fp_step_function:nnnN` { $\langle initial\ value \rangle$ } { $\langle step \rangle$ } { $\langle final\ value \rangle$ } $\langle function \rangle$

This function first evaluates the $\langle initial\ value \rangle$, $\langle step \rangle$ and $\langle final\ value \rangle$, each of which should be a floating point expression evaluating to a floating point number, not a tuple. The $\langle function \rangle$ is then placed in front of each $\langle value \rangle$ from the $\langle initial\ value \rangle$ to the $\langle final\ value \rangle$ in turn (using $\langle step \rangle$ between each $\langle value \rangle$). The $\langle step \rangle$ must be non-zero. If the $\langle step \rangle$ is positive, the loop stops when the $\langle value \rangle$ becomes larger than the $\langle final\ value \rangle$. If the $\langle step \rangle$ is negative, the loop stops when the $\langle value \rangle$ becomes smaller than the $\langle final\ value \rangle$. The $\langle function \rangle$ should absorb one numerical argument. For example

```
\cs_set:Npn \my_func:n #1 { [I~saw~#1] \quad }
\fp_step_function:nnnN { 1.0 } { 0.1 } { 1.5 } \my_func:n
```

would print

```
[I saw 1.0]   [I saw 1.1]   [I saw 1.2]   [I saw 1.3]   [I saw 1.4]   [I saw 1.5]
```

TpXhackers note: Due to rounding, it may happen that adding the $\langle step \rangle$ to the $\langle value \rangle$ does not change the $\langle value \rangle$; such cases give an error, as they would otherwise lead to an infinite loop.

`\fp_step_inline:nnnn`

New: 2016-11-21
Updated: 2016-12-06

`\fp_step_inline:nnnn` { $\langle initial\ value \rangle$ } { $\langle step \rangle$ } { $\langle final\ value \rangle$ } { $\langle code \rangle$ }

This function first evaluates the $\langle initial\ value \rangle$, $\langle step \rangle$ and $\langle final\ value \rangle$, all of which should be floating point expressions evaluating to a floating point number, not a tuple. Then for each $\langle value \rangle$ from the $\langle initial\ value \rangle$ to the $\langle final\ value \rangle$ in turn (using $\langle step \rangle$ between each $\langle value \rangle$), the $\langle code \rangle$ is inserted into the input stream with `#1` replaced by the current $\langle value \rangle$. Thus the $\langle code \rangle$ should define a function of one argument (`#1`).

`\fp_step_variable:nnnNn`

New: 2017-04-12

`\fp_step_variable:nnnNn`
{ $\langle initial\ value \rangle$ } { $\langle step \rangle$ } { $\langle final\ value \rangle$ } $\langle tl\ var \rangle$ { $\langle code \rangle$ }

This function first evaluates the $\langle initial\ value \rangle$, $\langle step \rangle$ and $\langle final\ value \rangle$, all of which should be floating point expressions evaluating to a floating point number, not a tuple. Then for each $\langle value \rangle$ from the $\langle initial\ value \rangle$ to the $\langle final\ value \rangle$ in turn (using $\langle step \rangle$ between each $\langle value \rangle$), the $\langle code \rangle$ is inserted into the input stream, with the $\langle tl\ var \rangle$ defined as the current $\langle value \rangle$. Thus the $\langle code \rangle$ should make use of the $\langle tl\ var \rangle$.

6 Some useful constants, and scratch variables

`\c_zero_fp`
`\c_minus_zero_fp`

New: 2012-05-08

Zero, with either sign.

`\c_one_fp`

New: 2012-05-08

One as an fp: useful for comparisons in some places.

<hr/> <code>\c_inf_fp</code> <code>\c_minus_inf_fp</code> <hr/> New: 2012-05-08 <hr/>	<p>Infinity, with either sign. These can be input directly in a floating point expression as <code>inf</code> and <code>-inf</code>.</p>
<hr/> <code>\c_e_fp</code> <hr/> Updated: 2012-05-08 <hr/>	<p>The value of the base of the natural logarithm, $e = \exp(1)$.</p>
<hr/> <code>\c_pi_fp</code> <hr/> Updated: 2013-11-17 <hr/>	<p>The value of π. This can be input directly in a floating point expression as <code>pi</code>.</p>
<hr/> <code>\c_one_degree_fp</code> <hr/> New: 2012-05-08 Updated: 2013-11-17 <hr/>	<p>The value of 1° in radians. Multiply an angle given in degrees by this value to obtain a result in radians. Note that trigonometric functions expecting an argument in radians or in degrees are both available. Within floating point expressions, this can be accessed as <code>deg</code>.</p>
<hr/> <code>\l_tmpa_fp</code> <code>\l_tmpb_fp</code> <hr/>	<p>Scratch floating points for local assignment. These are never used by the kernel code, and so are safe for use with any L^AT_EX3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.</p>
<hr/> <code>\g_tmpa_fp</code> <code>\g_tmpb_fp</code> <hr/>	<p>Scratch floating points for global assignment. These are never used by the kernel code, and so are safe for use with any L^AT_EX3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.</p>

7 Floating point exceptions

The functions defined in this section are experimental, and their functionality may be altered or removed altogether.

“Exceptions” may occur when performing some floating point operations, such as $0 / 0$, or $10 ** 1e9999$. The relevant IEEE standard defines 5 types of exceptions, of which we implement 4.

- *Overflow* occurs whenever the result of an operation is too large to be represented as a normal floating point number. This results in $\pm\infty$.
- *Underflow* occurs whenever the result of an operation is too close to 0 to be represented as a normal floating point number. This results in ± 0 .
- *Invalid operation* occurs for operations with no defined outcome, for instance $0/0$ or $\sin(\infty)$, and results in a NaN. It also occurs for conversion functions whose target type does not have the appropriate infinite or NaN value (*e.g.*, `\fp_to_dim:n`).
- *Division by zero* occurs when dividing a non-zero number by 0, or when evaluating functions at poles, *e.g.*, $\ln(0)$ or $\cot(0)$. This results in $\pm\infty$.

(*not yet*) *Inexact* occurs whenever the result of a computation is not exact, in other words, almost always. At the moment, this exception is entirely ignored in L^AT_EX3.

To each exception we associate a “flag”: `fp_overflow`, `fp_underflow`, `fp_invalid_operation` and `fp_division_by_zero`. The state of these flags can be tested and modified with commands from `l3flag`

By default, the “invalid operation” exception triggers an (expandable) error, and raises the corresponding flag. Other exceptions raise the corresponding flag but do not trigger an error. The behaviour when an exception occurs can be modified (using `\fp_trap:nn`) to either produce an error and raise the flag, or only raise the flag, or do nothing at all.

<code>\fp_trap:nn</code>	<code>\fp_trap:nn {<exception>} {<trap type>}</code>
New: 2012-07-19 Updated: 2017-02-13	All occurrences of the <code><exception></code> (<code>overflow</code> , <code>underflow</code> , <code>invalid_operation</code> or <code>division_by_zero</code>) within the current group are treated as <code><trap type></code> , which can be <ul style="list-style-type: none"> • none: the <code><exception></code> will be entirely ignored, and leave no trace; • flag: the <code><exception></code> will turn the corresponding flag on when it occurs; • error: additionally, the <code><exception></code> will halt the T_EX run and display some information about the current operation in the terminal.

This function is experimental, and may be altered or removed.

`flag_fp_overflow`
`flag_fp_underflow`
`flag_fp_invalid_operation`
`flag_fp_division_by_zero`

Flags denoting the occurrence of various floating-point exceptions.

8 Viewing floating points

<code>\fp_show:N</code>	<code>\fp_show:N <fp var></code>
<code>\fp_show:c</code>	<code>\fp_show:n {<floating point expression>}</code>
<code>\fp_show:n</code>	Evaluates the <code><floating point expression></code> and displays the result in the terminal.

New: 2012-05-08
Updated: 2015-08-07

<code>\fp_log:N</code>	<code>\fp_log:N <fp var></code>
<code>\fp_log:c</code>	<code>\fp_log:n {<floating point expression>}</code>
<code>\fp_log:n</code>	Evaluates the <code><floating point expression></code> and writes the result in the log file.

New: 2014-08-22
Updated: 2015-08-07

9 Floating point expressions

9.1 Input of floating point numbers

We support four types of floating point numbers:

- $\pm m \cdot 10^n$, a floating point number, with integer $1 \leq m \leq 10^{16}$, and $-10000 \leq n \leq 10000$;
- ± 0 , zero, with a given sign;
- $\pm \infty$, infinity, with a given sign;
- NaN, is “not a number”, and can be either quiet or signalling (*not yet*: this distinction is currently unsupported);

Normal floating point numbers are stored in base 10, with up to 16 significant figures.

On input, a normal floating point number consists of:

- $\langle sign \rangle$: a possibly empty string of + and - characters;
- $\langle significand \rangle$: a non-empty string of digits together with zero or one dot;
- $\langle exponent \rangle$ optionally: the character e or E, followed by a possibly empty string of + and - tokens, and a non-empty string of digits.

The sign of the resulting number is + if $\langle sign \rangle$ contains an even number of -, and - otherwise, hence, an empty $\langle sign \rangle$ denotes a non-negative input. The stored significand is obtained from $\langle significand \rangle$ by omitting the decimal separator and leading zeros, and rounding to 16 significant digits, filling with trailing zeros if necessary. In particular, the value stored is exact if the input $\langle significand \rangle$ has at most 16 digits. The stored $\langle exponent \rangle$ is obtained by combining the input $\langle exponent \rangle$ (0 if absent) with a shift depending on the position of the significand and the number of leading zeros.

A special case arises if the resulting $\langle exponent \rangle$ is either too large or too small for the floating point number to be represented. This results either in an overflow (the number is then replaced by $\pm \infty$), or an underflow (resulting in ± 0).

The result is thus ± 0 if and only if $\langle significand \rangle$ contains no non-zero digit (*i.e.*, consists only in characters 0, and an optional period), or if there is an underflow. Note that a single dot is currently a valid floating point number, equal to +0, but that is not guaranteed to remain true.

The $\langle significand \rangle$ must be non-empty, so e1 and e-1 are not valid floating point numbers. Note that the latter could be mistaken with the difference of “e” and 1. To avoid confusions, the base of natural logarithms cannot be input as e and should be input as exp(1) or \c_e_fp (which is faster).

Special numbers are input as follows:

- inf represents $+\infty$, and can be preceded by any $\langle sign \rangle$, yielding $\pm \infty$ as appropriate.
- nan represents a (quiet) non-number. It can be preceded by any sign, but that sign is ignored.
- Any unrecognizable string triggers an error, and produces a NaN.
- Note that commands such as \infy, \pi, or \sin *do not* work in floating point expressions. They may silently be interpreted as completely unexpected numbers, because integer constants (allowed in expressions) are commonly stored as mathematical characters.

9.2 Precedence of operators

We list here all the operations supported in floating point expressions, in order of decreasing precedence: operations listed earlier bind more tightly than operations listed below them.

- Function calls (`sin`, `ln`, *etc*).
- Binary `**` and `^` (right associative).
- Unary `+`, `-`, `!`.
- Implicit multiplication by juxtaposition (`2pi`) when neither factor is in parentheses.
- Binary `*` and `/`, implicit multiplication by juxtaposition with parentheses (for instance `3(4+5)`).
- Binary `+` and `-`.
- Comparisons `>=`, `!=`, `<?`, *etc*.
- Logical `and`, denoted by `&&`.
- Logical `or`, denoted by `||`.
- Ternary operator `?:` (right associative).
- Comma (to build tuples).

The precedence of operations can be overridden using parentheses. In particular, the precedence of juxtaposition implies that

$$\begin{aligned}1/2\text{pi} &= 1/(2\pi), \\1/2\text{pi}(\text{pi} + \text{pi}) &= (2\pi)^{-1}(\pi + \pi) \simeq 1, \\ \text{sin}2\text{pi} &= \sin(2)\pi \neq 0, \\ 2^2\text{max}(3, 5) &= 2^2 \max(3, 5) = 20, \\ 1\text{in}/1\text{cm} &= (1\text{in})/(1\text{cm}) = 2.54.\end{aligned}$$

Functions are called on the value of their argument, contrarily to `TEX` macros.

9.3 Operations

We now present the various operations allowed in floating point expressions, from the lowest precedence to the highest. When used as a truth value, a floating point expression is `false` if it is ± 0 , and `true` otherwise, including when it is `NaN` or a tuple such as `(0, 0)`. Tuples are only supported to some extent by operations that work with truth values (`?:`, `||`, `&&`, `!`), by comparisons (`!<=>?`), and by `+`, `-`, `*`, `/`. Unless otherwise specified, providing a tuple as an argument of any other operation yields the “invalid operation” exception and a `NaN` result.

```
?: \fp_eval:n { <operand1> ? <operand2> : <operand3> }
```

The ternary operator `?:` results in $\langle operand_2 \rangle$ if $\langle operand_1 \rangle$ is true (not ± 0), and $\langle operand_3 \rangle$ if $\langle operand_1 \rangle$ is false (± 0). All three $\langle operands \rangle$ are evaluated in all cases; they may be tuples. The operator is right associative, hence

```
\fp_eval:n
{
  1 + 3 > 4 ? 1 :
  2 + 4 > 5 ? 2 :
  3 + 5 > 6 ? 3 : 4
}
```

first tests whether $1 + 3 > 4$; since this isn't true, the branch following `:` is taken, and $2 + 4 > 5$ is compared; since this is true, the branch before `:` is taken, and everything else is (evaluated then) ignored. That allows testing for various cases in a concise manner, with the drawback that all computations are made in all cases.

```
|| \fp_eval:n { <operand1> || <operand2> }
```

If $\langle operand_1 \rangle$ is true (not ± 0), use that value, otherwise the value of $\langle operand_2 \rangle$. Both $\langle operands \rangle$ are evaluated in all cases; they may be tuples. In $\langle operand_1 \rangle || \langle operand_2 \rangle || \dots || \langle operands_n \rangle$, the first true (nonzero) $\langle operand \rangle$ is used and if all are zero the last one (± 0) is used.

```
&& \fp_eval:n { <operand1> && <operand2> }
```

If $\langle operand_1 \rangle$ is false (equal to ± 0), use that value, otherwise the value of $\langle operand_2 \rangle$. Both $\langle operands \rangle$ are evaluated in all cases; they may be tuples. In $\langle operand_1 \rangle \&\& \langle operand_2 \rangle \&\& \dots \&\& \langle operands_n \rangle$, the first false (± 0) $\langle operand \rangle$ is used and if none is zero the last one is used.

```
< \fp_eval:n
= {
>   <operand1> <relation1>
?   ...
   <operand_N> <relation_N>
   <operand_{N+1}>
}
```

Updated: 2013-12-14

Each $\langle relation \rangle$ consists of a non-empty string of `<`, `=`, `>`, and `?`, optionally preceded by `!`, and may not start with `?`. This evaluates to $+1$ if all comparisons $\langle operand_i \rangle \langle relation_i \rangle \langle operand_{i+1} \rangle$ are true, and $+0$ otherwise. All $\langle operands \rangle$ are evaluated (once) in all cases. See `\fp_compare:nTF` for details.

```
+ \fp_eval:n { <operand1> + <operand2> }
- \fp_eval:n { <operand1> - <operand2> }
```

Computes the sum or the difference of its two $\langle operands \rangle$. The “invalid operation” exception occurs for $\infty - \infty$. “Underflow” and “overflow” occur when appropriate. These operations supports the itemwise addition or subtraction of two tuples, but if they have a different number of items the “invalid operation” exception occurs and the result is NaN.

```

* \fp_eval:n { <operand1> * <operand2> }
/ \fp_eval:n { <operand1> / <operand2> }

```

Computes the product or the ratio of its two $\langle \text{operands} \rangle$. The “invalid operation” exception occurs for ∞/∞ , $0/0$, or $0 * \infty$. “Division by zero” occurs when dividing a finite non-zero number by ± 0 . “Underflow” and “overflow” occur when appropriate. When $\langle \text{operand}_1 \rangle$ is a tuple and $\langle \text{operand}_2 \rangle$ is a floating point number, each item of $\langle \text{operand}_1 \rangle$ is multiplied or divided by $\langle \text{operand}_2 \rangle$. Multiplication also supports the case where $\langle \text{operand}_1 \rangle$ is a floating point number and $\langle \text{operand}_2 \rangle$ a tuple. Other combinations yield an “invalid operation” exception and a NaN result.

```

+ \fp_eval:n { + <operand> }
- \fp_eval:n { - <operand> }
! \fp_eval:n { ! <operand> }

```

The unary $+$ does nothing, the unary $-$ changes the sign of the $\langle \text{operand} \rangle$ (for a tuple, of all its components), and $!$ $\langle \text{operand} \rangle$ evaluates to 1 if $\langle \text{operand} \rangle$ is false (is ± 0) and 0 otherwise (this is the `not` boolean function). Those operations never raise exceptions.

```

** \fp_eval:n { <operand1> ** <operand2> }
^ \fp_eval:n { <operand1> ^ <operand2> }

```

Raises $\langle \text{operand}_1 \rangle$ to the power $\langle \text{operand}_2 \rangle$. This operation is right associative, hence 2^{2^3} equals $2^{2^3} = 256$. If $\langle \text{operand}_1 \rangle$ is negative or -0 then: the result’s sign is $+$ if the $\langle \text{operand}_2 \rangle$ is infinite and $(-1)^p$ if the $\langle \text{operand}_2 \rangle$ is $p/5^q$ with p, q integers; the result is $+0$ if $\text{abs}(\langle \text{operand}_1 \rangle)^{\langle \text{operand}_2 \rangle}$ evaluates to zero; in other cases the “invalid operation” exception occurs because the sign cannot be determined. “Division by zero” occurs when raising ± 0 to a finite strictly negative power. “Underflow” and “overflow” occur when appropriate. If either operand is a tuple, “invalid operation” occurs.

```

abs \fp_eval:n { abs( <fpexpr> ) }

```

Computes the absolute value of the $\langle \text{fpexpr} \rangle$. If the operand is a tuple, “invalid operation” occurs. This operation does not raise exceptions in other cases. See also `\fp_abs:n`.

```

exp \fp_eval:n { exp( <fpexpr> ) }

```

Computes the exponential of the $\langle \text{fpexpr} \rangle$. “Underflow” and “overflow” occur when appropriate. If the operand is a tuple, “invalid operation” occurs.

```

fact \fp_eval:n { fact( <fpexpr> ) }

```

Computes the factorial of the $\langle \text{fpexpr} \rangle$. If the $\langle \text{fpexpr} \rangle$ is an integer between -0 and 3248 included, the result is finite and correctly rounded. Larger positive integers give $+\infty$ with “overflow”, while $\text{fact}(+\infty) = +\infty$ and $\text{fact}(\text{nan}) = \text{nan}$ with no exception. All other inputs give NaN with the “invalid operation” exception.

```

ln \fp_eval:n { ln( <fpexpr> ) }

```

Computes the natural logarithm of the $\langle \text{fpexpr} \rangle$. Negative numbers have no (real) logarithm, hence the “invalid operation” is raised in that case, including for $\ln(-0)$. “Division by zero” occurs when evaluating $\ln(+0) = -\infty$. “Underflow” and “overflow” occur when appropriate. If the operand is a tuple, “invalid operation” occurs.

<hr/> logb <hr/>	★	<code>\fp_eval:n { logb(<fpexpr>) }</code>	
<hr/> New: 2018-11-03 <hr/>			Determines the exponent of the $\langle fpexpr \rangle$, namely the floor of the base-10 logarithm of its absolute value. “Division by zero” occurs when evaluating $\logb(\pm 0) = -\infty$. Other special values are $\logb(\pm\infty) = +\infty$ and $\logb(\text{NaN}) = \text{NaN}$. If the operand is a tuple or is NaN, then “invalid operation” occurs and the result is NaN.
<hr/> max min <hr/>		<code>\fp_eval:n { max(<fpexpr₁> , <fpexpr₂> , ...) }</code> <code>\fp_eval:n { min(<fpexpr₁> , <fpexpr₂> , ...) }</code>	Evaluates each $\langle fpexpr \rangle$ and computes the largest (smallest) of those. If any of the $\langle fpexpr \rangle$ is a NaN or tuple, the result is NaN. If any operand is a tuple, “invalid operation” occurs; these operations do not raise exceptions in other cases.
<hr/> round trunc ceil floor <hr/>		<code>\fp_eval:n { round (<fpexpr>) }</code> <code>\fp_eval:n { round (<fpexpr₁> , <fpexpr₂>) }</code> <code>\fp_eval:n { round (<fpexpr₁> , <fpexpr₂> , <fpexpr₃>) }</code>	Only round accepts a third argument. Evaluates $\langle fpexpr_1 \rangle = x$ and $\langle fpexpr_2 \rangle = n$ and $\langle fpexpr_3 \rangle = t$ then rounds x to n places. If n is an integer, this rounds x to a multiple of 10^{-n} ; if $n = +\infty$, this always yields x ; if $n = -\infty$, this yields one of ± 0 , $\pm\infty$, or NaN; if $n = \text{NaN}$, this yields NaN; if n is neither $\pm\infty$ nor an integer, then an “invalid operation” exception is raised. When $\langle fpexpr_2 \rangle$ is omitted, $n = 0$, <i>i.e.</i> , $\langle fpexpr_1 \rangle$ is rounded to an integer. The rounding direction depends on the function. <ul style="list-style-type: none"> • round yields the multiple of 10^{-n} closest to x, with ties (x half-way between two such multiples) rounded as follows. If t is nan (or not given) the even multiple is chosen (“ties to even”), if $t = \pm 0$ the multiple closest to 0 is chosen (“ties to zero”), if t is positive/negative the multiple closest to $\infty/-\infty$ is chosen (“ties towards positive/negative infinity”). • floor yields the largest multiple of 10^{-n} smaller or equal to x (“round towards negative infinity”); • ceil yields the smallest multiple of 10^{-n} greater or equal to x (“round towards positive infinity”); • trunc yields a multiple of 10^{-n} with the same sign as x and with the largest absolute value less than that of x (“round towards zero”). <p>“Overflow” occurs if x is finite and the result is infinite (this can only happen if $\langle fpexpr_2 \rangle < -9984$). If any operand is a tuple, “invalid operation” occurs.</p>
<hr/> New: 2013-12-14 Updated: 2015-08-08 <hr/>			
<hr/> sign <hr/>		<code>\fp_eval:n { sign(<fpexpr>) }</code>	Evaluates the $\langle fpexpr \rangle$ and determines its sign: +1 for positive numbers and for $+\infty$, -1 for negative numbers and for $-\infty$, ± 0 for ± 0 , and NaN for NaN. If the operand is a tuple, “invalid operation” occurs. This operation does not raise exceptions in other cases.

<code>sin</code>	<code>\fp_eval:n { sin(<fpexpr>) }</code>
<code>cos</code>	<code>\fp_eval:n { cos(<fpexpr>) }</code>
<code>tan</code>	<code>\fp_eval:n { tan(<fpexpr>) }</code>
<code>cot</code>	<code>\fp_eval:n { cot(<fpexpr>) }</code>
<code>csc</code>	<code>\fp_eval:n { csc(<fpexpr>) }</code>
<code>sec</code>	<code>\fp_eval:n { sec(<fpexpr>) }</code>

Updated: 2013-11-17

Computes the sine, cosine, tangent, cotangent, cosecant, or secant of the $\langle fpexpr \rangle$ given in radians. For arguments given in degrees, see `sind`, `cosd`, *etc.* Note that since π is irrational, `sin(8pi)` is not quite zero, while its analogue `sind(8 × 180)` is exactly zero. The trigonometric functions are undefined for an argument of $\pm\infty$, leading to the “invalid operation” exception. Additionally, evaluating tangent, cotangent, cosecant, or secant at one of their poles leads to a “division by zero” exception. “Underflow” and “overflow” occur when appropriate. If the operand is a tuple, “invalid operation” occurs.

<code>sind</code>	<code>\fp_eval:n { sind(<fpexpr>) }</code>
<code>cosd</code>	<code>\fp_eval:n { cosd(<fpexpr>) }</code>
<code>tand</code>	<code>\fp_eval:n { tand(<fpexpr>) }</code>
<code>cotd</code>	<code>\fp_eval:n { cotd(<fpexpr>) }</code>
<code>cscd</code>	<code>\fp_eval:n { cscd(<fpexpr>) }</code>
<code>secd</code>	<code>\fp_eval:n { secd(<fpexpr>) }</code>

New: 2013-11-02

Computes the sine, cosine, tangent, cotangent, cosecant, or secant of the $\langle fpexpr \rangle$ given in degrees. For arguments given in radians, see `sin`, `cos`, *etc.* Note that since π is irrational, `sin(8pi)` is not quite zero, while its analogue `sind(8 × 180)` is exactly zero. The trigonometric functions are undefined for an argument of $\pm\infty$, leading to the “invalid operation” exception. Additionally, evaluating tangent, cotangent, cosecant, or secant at one of their poles leads to a “division by zero” exception. “Underflow” and “overflow” occur when appropriate. If the operand is a tuple, “invalid operation” occurs.

<code>asin</code>	<code>\fp_eval:n { asin(<fpexpr>) }</code>
<code>acos</code>	<code>\fp_eval:n { acos(<fpexpr>) }</code>
<code>acsc</code>	<code>\fp_eval:n { acsc(<fpexpr>) }</code>
<code>asec</code>	<code>\fp_eval:n { asec(<fpexpr>) }</code>

New: 2013-11-02

Computes the arcsine, arccosine, arccosecant, or arcsecant of the $\langle fpexpr \rangle$ and returns the result in radians, in the range $[-\pi/2, \pi/2]$ for `asin` and `acsc` and $[0, \pi]$ for `acos` and `asec`. For a result in degrees, use `asind`, *etc.* If the argument of `asin` or `acos` lies outside the range $[-1, 1]$, or the argument of `acsc` or `asec` inside the range $(-1, 1)$, an “invalid operation” exception is raised. “Underflow” and “overflow” occur when appropriate. If the operand is a tuple, “invalid operation” occurs.

<code>asind</code>	<code>\fp_eval:n { asind(<fpexpr>) }</code>
<code>acosd</code>	<code>\fp_eval:n { acosd(<fpexpr>) }</code>
<code>acscd</code>	<code>\fp_eval:n { acscd(<fpexpr>) }</code>
<code>asecd</code>	<code>\fp_eval:n { asecd(<fpexpr>) }</code>

New: 2013-11-02

Computes the arcsine, arccosine, arccosecant, or arcsecant of the $\langle fpexpr \rangle$ and returns the result in degrees, in the range $[-90, 90]$ for `asin` and `acsc` and $[0, 180]$ for `acos` and `asec`. For a result in radians, use `asin`, *etc.* If the argument of `asin` or `acos` lies outside the range $[-1, 1]$, or the argument of `acsc` or `asec` inside the range $(-1, 1)$, an “invalid operation” exception is raised. “Underflow” and “overflow” occur when appropriate. If the operand is a tuple, “invalid operation” occurs.

atan	<code>\fp_eval:n { atan(<fpexpr>) }</code>
acot	<code>\fp_eval:n { atan(<fpexpr₁> , <fpexpr₂>) }</code>
<hr/>	
New: 2013-11-02	<code>\fp_eval:n { acot(<fpexpr>) }</code>
	<code>\fp_eval:n { acot(<fpexpr₁> , <fpexpr₂>) }</code>

Those functions yield an angle in radians: **atand** and **acotd** are their analogs in degrees. The one-argument versions compute the arctangent or arccotangent of the $\langle fpexpr \rangle$: arctangent takes values in the range $[-\pi/2, \pi/2]$, and arccotangent in the range $[0, \pi]$. The two-argument arctangent computes the angle in polar coordinates of the point with Cartesian coordinates $(\langle fpexpr_2 \rangle, \langle fpexpr_1 \rangle)$: this is the arctangent of $\langle fpexpr_1 \rangle / \langle fpexpr_2 \rangle$, possibly shifted by π depending on the signs of $\langle fpexpr_1 \rangle$ and $\langle fpexpr_2 \rangle$. The two-argument arccotangent computes the angle in polar coordinates of the point $(\langle fpexpr_1 \rangle, \langle fpexpr_2 \rangle)$, equal to the arccotangent of $\langle fpexpr_1 \rangle / \langle fpexpr_2 \rangle$, possibly shifted by π . Both two-argument functions take values in the wider range $[-\pi, \pi]$. The ratio $\langle fpexpr_1 \rangle / \langle fpexpr_2 \rangle$ need not be defined for the two-argument arctangent: when both expressions yield ± 0 , or when both yield $\pm \infty$, the resulting angle is one of $\{\pm\pi/4, \pm 3\pi/4\}$ depending on signs. The “underflow” exception can occur. If any operand is a tuple, “invalid operation” occurs.

atand	<code>\fp_eval:n { atand(<fpexpr>) }</code>
acotd	<code>\fp_eval:n { atand(<fpexpr₁> , <fpexpr₂>) }</code>
<hr/>	
New: 2013-11-02	<code>\fp_eval:n { acotd(<fpexpr>) }</code>
	<code>\fp_eval:n { acotd(<fpexpr₁> , <fpexpr₂>) }</code>

Those functions yield an angle in degrees: **atand** and **acotd** are their analogs in radians. The one-argument versions compute the arctangent or arccotangent of the $\langle fpexpr \rangle$: arctangent takes values in the range $[-90, 90]$, and arccotangent in the range $[0, 180]$. The two-argument arctangent computes the angle in polar coordinates of the point with Cartesian coordinates $(\langle fpexpr_2 \rangle, \langle fpexpr_1 \rangle)$: this is the arctangent of $\langle fpexpr_1 \rangle / \langle fpexpr_2 \rangle$, possibly shifted by 180 depending on the signs of $\langle fpexpr_1 \rangle$ and $\langle fpexpr_2 \rangle$. The two-argument arccotangent computes the angle in polar coordinates of the point $(\langle fpexpr_1 \rangle, \langle fpexpr_2 \rangle)$, equal to the arccotangent of $\langle fpexpr_1 \rangle / \langle fpexpr_2 \rangle$, possibly shifted by 180. Both two-argument functions take values in the wider range $[-180, 180]$. The ratio $\langle fpexpr_1 \rangle / \langle fpexpr_2 \rangle$ need not be defined for the two-argument arctangent: when both expressions yield ± 0 , or when both yield $\pm \infty$, the resulting angle is one of $\{\pm 45, \pm 135\}$ depending on signs. The “underflow” exception can occur. If any operand is a tuple, “invalid operation” occurs.

sqrt	<code>\fp_eval:n { sqrt(<fpexpr>) }</code>
-------------	--

New: 2013-12-14 Computes the square root of the $\langle fpexpr \rangle$. The “invalid operation” is raised when the $\langle fpexpr \rangle$ is negative or is a tuple; no other exception can occur. Special values yield $\sqrt{-0} = -0$, $\sqrt{+0} = +0$, $\sqrt{+\infty} = +\infty$ and $\sqrt{\text{NaN}} = \text{NaN}$.

<hr/> rand <hr/>	<code>\fp_eval:n { rand() }</code>
<hr/> <small>New: 2016-12-05</small> <hr/>	<p>Produces a pseudo-random floating-point number (multiple of 10^{-16}) between 0 included and 1 excluded. This is not available in older versions of $\text{X}\text{_}\text{T}\text{E}\text{X}$. The random seed can be queried using <code>\sys_rand_seed:</code> and set using <code>\sys_gset_rand_seed:n</code>.</p> <p>TEXhackers note: This is based on pseudo-random numbers provided by the engine’s primitive <code>\pdfuniformdeviate</code> in $\text{pdf}\text{_}\text{T}\text{E}\text{X}$, $\text{p}\text{T}\text{E}\text{X}$, $\text{up}\text{T}\text{E}\text{X}$ and <code>\uniformdeviate</code> in $\text{Lua}\text{_}\text{T}\text{E}\text{X}$ and $\text{X}\text{_}\text{T}\text{E}\text{X}$. The underlying code is based on Metapost, which follows an additive scheme recommended in Section 3.6 of “The Art of Computer Programming, Volume 2”.</p> <p>While we are more careful than <code>\uniformdeviate</code> to preserve uniformity of the underlying stream of 28-bit pseudo-random integers, these pseudo-random numbers should of course not be relied upon for serious numerical computations nor cryptography.</p>
<hr/> randint <hr/>	<code>\fp_eval:n { randint(<fpexpr>) }</code>
<hr/> <small>New: 2016-12-05</small> <hr/>	<code>\fp_eval:n { randint(<fpexpr₁> , <fpexpr₂>) }</code> <p>Produces a pseudo-random integer between 1 and $\langle fpexpr \rangle$ or between $\langle fpexpr_1 \rangle$ and $\langle fpexpr_2 \rangle$ inclusive. The bounds must be integers in the range $(-10^{16}, 10^{16})$ and the first must be smaller or equal to the second. See rand for important comments on how these pseudo-random numbers are generated.</p>
<hr/> inf nan <hr/>	The special values $+\infty$, $-\infty$, and NaN are represented as inf , -inf and nan (see <code>\c_minus_inf_fp</code> , <code>\c_minus_inf_fp</code> and <code>\c_nan_fp</code>).
<hr/> pi <hr/>	The value of π (see <code>\c_pi_fp</code>).
<hr/> deg <hr/>	The value of 1° in radians (see <code>\c_one_degree_fp</code>).

<hr/>	Those units of measurement are equal to their values in <code>pt</code> , namely
<code>em</code>	
<code>ex</code>	
<code>in</code>	$1\text{ in} = 72.27\text{ pt}$
<code>pt</code>	$1\text{ pt} = 1\text{ pt}$
<code>pc</code>	
<code>cm</code>	$1\text{ pc} = 12\text{ pt}$
<code>mm</code>	
<code>dd</code>	$1\text{ cm} = \frac{1}{2.54}\text{ in} = 28.45275590551181\text{ pt}$
<code>cc</code>	
<code>nd</code>	$1\text{ mm} = \frac{1}{25.4}\text{ in} = 2.845275590551181\text{ pt}$
<code>nc</code>	
<code>bp</code>	$1\text{ dd} = 0.376065\text{ mm} = 1.07000856496063\text{ pt}$
<code>sp</code>	$1\text{ cc} = 12\text{ dd} = 12.84010277952756\text{ pt}$
<hr/>	$1\text{ nd} = 0.375\text{ mm} = 1.066978346456693\text{ pt}$
	$1\text{ nc} = 12\text{ nd} = 12.80374015748031\text{ pt}$
	$1\text{ bp} = \frac{1}{72}\text{ in} = 1.00375\text{ pt}$
	$1\text{ sp} = 2^{-16}\text{ pt} = 1.52587890625 \times 10^{-5}\text{ pt}.$

The values of the (font-dependent) units `em` and `ex` are gathered from \TeX when the surrounding floating point expression is evaluated.

<hr/>	Other names for 1 and $+0$.
<code>true</code>	
<code>false</code>	
<hr/>	

<hr/>	
<code>\fp_abs:n</code> *	<code>\fp_abs:n {⟨floating point expression⟩}</code>
New: 2012-05-14	Evaluates the <i>⟨floating point expression⟩</i> as described for <code>\fp_eval:n</code> and leaves the absolute value of the result in the input stream. If the argument is $\pm\infty$, <code>NaN</code> or a tuple, “invalid operation” occurs. Within floating point expressions, <code>abs()</code> can be used; it accepts $\pm\infty$ and <code>NaN</code> as arguments.
Updated: 2012-07-08	
<hr/>	

<hr/>	
<code>\fp_max:nn</code> *	<code>\fp_max:nn {⟨fp expression 1⟩} {⟨fp expression 2⟩}</code>
<code>\fp_min:nn</code> *	Evaluates the <i>⟨floating point expressions⟩</i> as described for <code>\fp_eval:n</code> and leaves the resulting larger (<code>max</code>) or smaller (<code>min</code>) value in the input stream. If the argument is a tuple, “invalid operation” occurs, but no other case raises exceptions. Within floating point expressions, <code>max()</code> and <code>min()</code> can be used.
New: 2012-09-26	
<hr/>	

10 Disclaimer and roadmap

The package may break down if the escape character is among `0123456789_+`, or if it receives a \TeX primitive conditional affected by `\exp_not:N`.

The following need to be done. I’ll try to time-order the items.

- Function to count items in a tuple (and to determine if something is a tuple).
- Decide what exponent range to consider.

- Support signalling `nan`.
- Modulo and remainder, and rounding function `quantize` (and its friends analogous to `trunc`, `ceil`, `floor`).
- `\fp_format:nn {<fpepr>} {<format>}`, but what should `<format>` be? More general pretty printing?
- Add `and`, `or`, `xor`? Perhaps under the names `all`, `any`, and `xor`?
- Add `log(x,b)` for logarithm of x in base b .
- `hypot` (Euclidean length). Cartesian-to-polar transform.
- Hyperbolic functions `cosh`, `sinh`, `tanh`.
- Inverse hyperbolics.
- Base conversion, input such as `0xAB.CDEF`.
- Factorial (not with `!`), gamma function.
- Improve coefficients of the `sin` and `tan` series.
- Treat upper and lower case letters identically in identifiers, and ignore underscores.
- Add an `array(1,2,3)` and `i=complex(0,1)`.
- Provide an experimental `map` function? Perhaps easier to implement if it is a single character, `@sin(1,2)`?
- Provide an `isnan` function analogue of `\fp_if_nan:nTF`?
- Support keyword arguments?

`Pgfmath` also provides box-measurements (depth, height, width), but boxes are not possible expandably.

Bugs, and tests to add.

- Check that functions are monotonic when they should.
- Add exceptions to `?:`, `!<=>?`, `&&`, `||`, and `!`.
- Logarithms of numbers very close to 1 are inaccurate.
- When rounding towards $-\infty$, `\dim_to_fp:n {Opt}` should return -0 , not $+0$.
- The result of $(\pm 0) + (\pm 0)$, of $x + (-x)$, and of $(-x) + x$ should depend on the rounding mode.
- `0e9999999999` gives a \TeX “number too large” error.
- Subnormals are not implemented.

Possible optimizations/improvements.

- Document that `l3trial/l3fp-types` introduces tools for adding new types.
- In subsection 9.1, write a grammar.

- It would be nice if the `parse` auxiliaries for each operation were set up in the corresponding module, rather than centralizing in `l3fp-parse`.
- Some functions should get an `_o` ending to indicate that they expand after their result.
- More care should be given to distinguish expandable/restricted expandable (auxiliary and internal) functions.
- The code for the `ternary` set of functions is ugly.
- There are many `~` missing in the doc to avoid bad line-breaks.
- The algorithm for computing the logarithm of the significand could be made to use a 5 terms Taylor series instead of 10 terms by taking $c = 2000/([200x]+1) \in [10, 95]$ instead of $c \in [1, 10]$. Also, it would then be possible to simplify the computation of t . However, we would then have to hard-code the logarithms of 44 small integers instead of 9.
- Improve notations in the explanations of the division algorithm (`l3fp-basics`).
- Understand and document `_fp_basics_pack_weird_low:NNNNw` and `_fp_basics_pack_weird_high:NNNNNNNNw` better. Move the other `basics_pack` auxiliaries to `l3fp-aux` under a better name.
- Find out if underflow can really occur for trigonometric functions, and redoc as appropriate.
- Add bibliography. Some of Kahan’s articles, some previous T_EX fp packages, the international standards,...
- Also take into account the “inexact” exception?
- Support multi-character prefix operators (*e.g.*, `@/` or whatever)?

Part XXIV

The l3farray package: fast global floating point arrays

1 l3farray documentation

For applications requiring heavy use of floating points, this module provides arrays which can be accessed in constant time (contrast l3seq, where access time is linear). The interface is very close to that of l3intarray. The size of the array is fixed and must be given at point of initialisation

<code>\farray_new:Nn</code>	<code>\farray_new:Nn <farray var> {<size>}</code>
-----------------------------	---

New: 2018-05-05

Evaluates the integer expression *<size>* and allocates an *<floating point array variable>* with that number of (zero) entries. The variable name should start with `\g_` because assignments are always global.

<code>\farray_count:N</code> ★	<code>\farray_count:N <farray var></code>
--------------------------------	---

New: 2018-05-05

Expands to the number of entries in the *<floating point array variable>*. This is performed in constant time.

<code>\farray_gset:Nnn</code>	<code>\farray_gset:Nnn <farray var> {<position>} {<value>}</code>
-------------------------------	---

New: 2018-05-05

Stores the result of evaluating the floating point expression *<value>* into the *<floating point array variable>* at the (integer expression) *<position>*. If the *<position>* is not between 1 and the `\farray_count:N`, an error occurs. Assignments are always global.

<code>\farray_gzero:N</code>	<code>\farray_gzero:N <farray var></code>
------------------------------	---

New: 2018-05-05

Sets all entries of the *<floating point array variable>* to +0. Assignments are always global.

<code>\farray_item:Nn</code> ★	<code>\farray_item:Nn <farray var> {<position>}</code>
--------------------------------	--

`\farray_item_to_tl:Nn` ★

New: 2018-05-05

Applies `\fp_use:N` or `\fp_to_tl:N` (respectively) to the floating point entry stored at the (integer expression) *<position>* in the *<floating point array variable>*. If the *<position>* is not between 1 and the `\farray_count:N`, an error occurs.

Part XXV

The l3cctab package

Category code tables

A category code table enables rapid switching of all category codes in one operation. For LuaTeX, this is possible over the entire Unicode range. For other engines, only the 8-bit range (0–255) is covered by such tables.

1 Creating and initialising category code tables

<hr/> <code>\cctab_new:N</code> <code>\cctab_new:c</code> <hr/> <small>Updated: 2020-07-02</small> <hr/>	<code>\cctab_new:N</code> \langle <i>category code table</i> \rangle Creates a new \langle <i>category code table</i> \rangle variable or raises an error if the name is already taken. The declaration is global. The \langle <i>category code table</i> \rangle is initialised with the codes as used by <code>iniTeX</code> .
<hr/> <code>\cctab_const:Nn</code> <code>\cctab_const:cn</code> <hr/> <small>Updated: 2020-07-07</small> <hr/>	<code>\cctab_const:Nn</code> \langle <i>category code table</i> \rangle $\{ \langle$ <i>category code set up</i> $\rangle \}$ Creates a new \langle <i>category code table</i> \rangle , applies (in a group) the \langle <i>category code set up</i> \rangle on top of <code>iniTeX</code> settings, then saves them globally as a constant table. The \langle <i>category code set up</i> \rangle can include a call to <code>\cctab_select:N</code> .
<hr/> <code>\cctab_gset:Nn</code> <code>\cctab_gset:cn</code> <hr/> <small>Updated: 2020-07-07</small> <hr/>	<code>\cctab_gset:Nn</code> \langle <i>category code table</i> \rangle $\{ \langle$ <i>category code set up</i> $\rangle \}$ Starting from the <code>iniTeX</code> category codes, applies (in a group) the \langle <i>category code set up</i> \rangle , then saves them globally in the \langle <i>category code table</i> \rangle . The \langle <i>category code set up</i> \rangle can include a call to <code>\cctab_select:N</code> .

2 Using category code tables

<hr/> <code>\cctab_begin:N</code> <code>\cctab_begin:c</code> <hr/> <small>Updated: 2020-07-02</small> <hr/>	<code>\cctab_begin:N</code> \langle <i>category code table</i> \rangle Switches locally the category codes in force to those stored in the \langle <i>category code table</i> \rangle . The prevailing codes before the function is called are added to a stack, for use with <code>\cctab_end:.</code> This function does not start a TeX group.
<hr/> <code>\cctab_end:</code> <hr/> <small>Updated: 2020-07-02</small> <hr/>	<code>\cctab_end:</code> Ends the scope of a \langle <i>category code table</i> \rangle started using <code>\cctab_begin:N</code> , returning the codes to those in force before the matching <code>\cctab_begin:N</code> was used. This must be used within the same TeX group (and at the same TeX group level) as the matching <code>\cctab_begin:N</code> .
<hr/> <code>\cctab_select:N</code> <hr/> <small>New: 2020-05-19 Updated: 2020-07-02</small> <hr/>	<code>\cctab_select:N</code> \langle <i>category code table</i> \rangle Selects the \langle <i>category code table</i> \rangle for the scope of the current group. This is in particular useful in the \langle <i>setup</i> \rangle arguments of <code>\tl_set_rescan:Nnn</code> , <code>\tl_rescan:nn</code> , <code>\cctab_const:Nn</code> , and <code>\cctab_gset:Nn</code> .

3 Category code table conditionals

<hr/> <code>\cctab_if_exist_p:N</code> *	<code>\cctab_if_exist_p:N</code> \langle <i>category code table</i> \rangle
<code>\cctab_if_exist_p:c</code> *	<code>\cctab_if_exist:NTF</code> \langle <i>category code table</i> \rangle $\{\langle$ <i>true code</i> $\rangle\}$ $\{\langle$ <i>false code</i> $\rangle\}$
<code>\cctab_if_exist:NTF</code> *	Tests whether the \langle <i>category code table</i> \rangle is currently defined. This does not check that the \langle <i>category code table</i> \rangle really is a category code table.
<code>\cctab_if_exist:cTF</code> *	

4 Constant category code tables

<hr/> <code>\c_code_cctab</code>	Category code table for the <code>expl3</code> code environment; this does <i>not</i> include <code>@</code> , which is retained as an “other” character.
Updated: 2020-07-10	

<hr/> <code>\c_document_cctab</code>	Category code table for a standard L ^A T _E X document, as set by the L ^A T _E X kernel. In particular, the upper-half of the 8-bit range will be set to “active” with pdfT _E X <i>only</i> . No <code>babel</code> shorthands will be activated.
Updated: 2020-07-08	

<hr/> <code>\c_initex_cctab</code>	Category code table as set up by <code>iniT_EX</code> .
Updated: 2020-07-02	

<hr/> <code>\c_other_cctab</code>	Category code table where all characters have category code 12 (other).
Updated: 2020-07-02	

<hr/> <code>\c_str_cctab</code>	Category code table where all characters have category code 12 (other) with the exception of spaces, which have category code 10 (space).
Updated: 2020-07-02	

Part XXVI

The l3sort package

Sorting functions

1 Controlling sorting

L^AT_EX3 comes with a facility to sort list variables (sequences, token lists, or comma-lists) according to some user-defined comparison. For instance,

```
\clist_set:Nn \l_foo_clist { 3 , 01 , -2 , 5 , +1 }
\clist_sort:Nn \l_foo_clist
{
  \int_compare:nNnTF { #1 } > { #2 }
  { \sort_return_swapped: }
  { \sort_return_same: }
}
```

results in `\l_foo_clist` holding the values `{ -2 , 01 , +1 , 3 , 5 }` sorted in non-decreasing order.

The code defining the comparison should call `\sort_return_swapped:` if the two items given as `#1` and `#2` are not in the correct order, and otherwise it should call `\sort_return_same:` to indicate that the order of this pair of items should not be changed.

For instance, a *comparison code* consisting only of `\sort_return_same:` with no test yields a trivial sort: the final order is identical to the original order. Conversely, using a *comparison code* consisting only of `\sort_return_swapped:` reverses the list (in a fairly inefficient way).

T_EXhackers note: The current implementation is limited to sorting approximately 20000 items (40000 in LuaT_EX), depending on what other packages are loaded.

Internally, the code from l3sort stores items in `\toks` registers allocated locally. Thus, the *comparison code* should not call `\newtoks` or other commands that allocate new `\toks` registers. On the other hand, altering the value of a previously allocated `\toks` register is not a problem.

```
\sort_return_same:
\sort_return_swapped:
```

New: 2017-02-06

```
\seq_sort:Nn <seq var>
{ ... \sort_return_same: or \sort_return_swapped: ... }
```

Indicates whether to keep the order or swap the order of two items that are compared in the sorting code. Only one of the `\sort_return_...` functions should be used by the code, according to the results of some tests on the items `#1` and `#2` to be compared.

Part XXVII

The l3tl-analysis package: Analysing token lists

1 l3tl-analysis documentation

This module provides functions that are particularly useful in the l3regex module for mapping through a token list one $\langle token \rangle$ at a time (including begin-group/end-group tokens). For `\tl_analysis_map_inline:Nn` or `\tl_analysis_map_inline:nn`, the token list is given as an argument; the analogous function `\peek_analysis_map_inline:n` documented in l3token finds tokens in the input stream instead. In both cases the user provides $\langle inline code \rangle$ that receives three arguments for each $\langle token \rangle$:

- $\langle tokens \rangle$, which both o-expand and x-expand to the $\langle token \rangle$. The detailed form of $\langle tokens \rangle$ may change in later releases.
- $\langle char code \rangle$, a decimal representation of the character code of the $\langle token \rangle$, -1 if it is a control sequence.
- $\langle catcode \rangle$, a capital hexadecimal digit which denotes the category code of the $\langle token \rangle$ (0: control sequence, 1: begin-group, 2: end-group, 3: math shift, 4: alignment tab, 6: parameter, 7: superscript, 8: subscript, A: space, B: letter, C: other, D: active). This can be converted to an integer by writing " $\langle catcode \rangle$ ".

In addition, there is a debugging function `\tl_analysis_show:n`, very similar to the `\ShowTokens` macro from the ted package.

```
\tl_analysis_show:N
\tl_analysis_show:n
```

New: 2018-04-09

```
\tl_analysis_show:n {\langle token list \rangle}
```

Displays to the terminal the detailed decomposition of the $\langle token list \rangle$ into tokens, showing the category code of each character token, the meaning of control sequences and active characters, and the value of registers.

```
\tl_analysis_map_inline:nn
\tl_analysis_map_inline:Nn
```

New: 2018-04-09

```
\tl_analysis_map_inline:nn {\langle token list \rangle} {\langle inline function \rangle}
```

Applies the $\langle inline function \rangle$ to each individual $\langle token \rangle$ in the $\langle token list \rangle$. The $\langle inline function \rangle$ receives three arguments as explained above. As all other mappings the mapping is done at the current group level, *i.e.* any local assignments made by the $\langle inline function \rangle$ remain in effect after the loop.

Part XXVIII

The `l3regex` package: Regular expressions in `TEX`

The `l3regex` package provides regular expression testing, extraction of submatches, splitting, and replacement, all acting on token lists. The syntax of regular expressions is mostly a subset of the PCRE syntax (and very close to POSIX), with some additions due to the fact that `TEX` manipulates tokens rather than characters. For performance reasons, only a limited set of features are implemented. Notably, back-references are not supported.

Let us give a few examples. After

```
\tl_set:Nn \l_my_tl { That~cat. }
\regex_replace_once:nnN { at } { is } \l_my_tl
```

the token list variable `\l_my_tl` holds the text “`This cat.`”, where the first occurrence of “`at`” was replaced by “`is`”. A more complicated example is a pattern to emphasize each word and add a comma after it:

```
\regex_replace_all:nnN { \w+ } { \c{emph}\cB\{ \0 \cE\} , } \l_my_tl
```

The `\w` sequence represents any “word” character, and `+` indicates that the `\w` sequence should be repeated as many times as possible (at least once), hence matching a word in the input token list. In the replacement text, `\0` denotes the full match (here, a word). The command `\emph` is inserted using `\c{emph}`, and its argument `\0` is put between braces `\cB\{` and `\cE\}`.

If a regular expression is to be used several times, it can be compiled once, and stored in a regex variable using `\regex_const:Nn`. For example,

```
\regex_const:Nn \c_foo_regex { \c{begin} \cB. (\c[~BE].*) \cE. }
```

stores in `\c_foo_regex` a regular expression which matches the starting marker for an environment: `\begin`, followed by a begin-group token (`\cB.`), then any number of tokens which are neither begin-group nor end-group character tokens (`\c[~BE].*`), ending with an end-group token (`\cE.`). As explained in the next section, the parentheses “capture” the result of `\c[~BE].*`, giving us access to the name of the environment when doing replacements.

1 Syntax of regular expressions

We start with a few examples, and encourage the reader to apply `\regex_show:n` to these regular expressions.

- `Cat` matches the word “Cat” capitalized in this way, but also matches the beginning of the word “Cattle”: use `\bCat\b` to match a complete word only.
- `[abc]` matches one letter among “a”, “b”, “c”; the pattern `(a|b|c)` matches the same three possible letters (but see the discussion of submatches below).
- `[A-Za-z]*` matches any number (due to the quantifier `*`) of Latin letters (not accented).

- `\c{[A-Za-z]*}` matches a control sequence made of Latin letters.
- `_[^_]*_` matches an underscore, any number of characters other than underscore, and another underscore; it is equivalent to `_.*?_` where `.` matches arbitrary characters and the lazy quantifier `*?` means to match as few characters as possible, thus avoiding matching underscores.
- `[\+|-]?[d+]` matches an explicit integer with at most one sign.
- `[\+|-_]*[d+_]*` matches an explicit integer with any number of `+` and `-` signs, with spaces allowed except within the mantissa, and surrounded by spaces.
- `[\+|-_]*(\d+|\d*\.\d+)_\d*` matches an explicit integer or decimal number; using `[.,]` instead of `\.` would allow the comma as a decimal marker.
- `[\+|-_]*(\d+|\d*\.\d+)_\d*((?i)pt|in|[cem]m|ex|[bs]p|[dn]d|[pcn]c)_\d*` matches an explicit dimension with any unit that T_EX knows, where `(?i)` means to treat lowercase and uppercase letters identically.
- `[\+|-_]*((?i)nan|inf|(\d+|\d*\.\d+)_\d*e[\+|-_]*/\d+)?_\d*` matches an explicit floating point number or the special values `nan` and `inf` (with signs and spaces allowed).
- `[\+|-_]*(\d+|\dC.)_\d*` matches an explicit integer or control sequence (without checking whether it is an integer variable).
- `\G.*?\K` at the beginning of a regular expression matches and discards (due to `\K`) everything between the end of the previous match (`\G`) and what is matched by the rest of the regular expression; this is useful in `\regex_replace_all:nnN` when the goal is to extract matches or submatches in a finer way than with `\regex_extract_all:nnN`.

While it is impossible for a regular expression to match only integer expressions, `[\+|-\(\)*\d+\)*(\+|-*/[\+|-\(\)*\d+\)\)]*` matches among other things all valid integer expressions (made only with explicit integers). One should follow it with further testing.

Most characters match exactly themselves, with an arbitrary category code. Some characters are special and must be escaped with a backslash (*e.g.*, `*` matches a star character). Some escape sequences of the form backslash-letter also have a special meaning (for instance `\d` matches any digit). As a rule,

- every alphanumeric character (`A-Z`, `a-z`, `0-9`) matches exactly itself, and should not be escaped, because `\A`, `\B`, ... have special meanings;
- non-alphanumeric printable ascii characters can (and should) always be escaped: many of them have special meanings (*e.g.*, use `\(`, `\)`, `\?`, `\.`);
- spaces should always be escaped (even in character classes);
- any other character may be escaped or not, without any effect: both versions match exactly that character.

Note that these rules play nicely with the fact that many non-alphanumeric characters are difficult to input into T_EX under normal category codes. For instance, `\abc%` matches the characters `\abc%` (with arbitrary category codes), but does not match the control sequence `\abc` followed by a percent character. Matching control sequences can be done using the `\c{<regex>}` syntax (see below).

Any special character which appears at a place where its special behaviour cannot apply matches itself instead (for instance, a quantifier appearing at the beginning of a string), after raising a warning.

Characters.

`\x{hh...}` Character with hex code `hh...`

`\xhh` Character with hex code `hh`.

`\a` Alarm (hex 07).

`\e` Escape (hex 1B).

`\f` Form-feed (hex 0C).

`\n` New line (hex 0A).

`\r` Carriage return (hex 0D).

`\t` Horizontal tab (hex 09).

Character types.

`.` A single period matches any token.

`\d` Any decimal digit.

`\h` Any horizontal space character, equivalent to `[\ \^^I]`: space and tab.

`\s` Any space character, equivalent to `[\ \^^I\^^J\^^L\^^M]`.

`\v` Any vertical space character, equivalent to `[\^^J\^^K\^^L\^^M]`. Note that `\^^K` is a vertical space, but not a space, for compatibility with Perl.

`\w` Any word character, *i.e.*, alphanumerics and underscore, equivalent to the explicit class `[A-Za-z0-9_]`.

`\D` Any token not matched by `\d`.

`\H` Any token not matched by `\h`.

`\N` Any token other than the `\n` character (hex 0A).

`\S` Any token not matched by `\s`.

`\V` Any token not matched by `\v`.

`\W` Any token not matched by `\w`.

Of those, `.`, `\D`, `\H`, `\N`, `\S`, `\V`, and `\W` match arbitrary control sequences.

Character classes match exactly one token in the subject.

`[...]` Positive character class. Matches any of the specified tokens.

[**^...**] Negative character class. Matches any token other than the specified characters.

x-y Within a character class, this denotes a range (can be used with escaped characters).

[:**<name>**:] Within a character class (one more set of brackets), this denotes the POSIX character class **<name>**, which can be **alnum**, **alpha**, **ascii**, **blank**, **cntrl**, **digit**, **graph**, **lower**, **print**, **punct**, **space**, **upper**, **word**, or **xdigit**.

[:**~<name>**:] Negative POSIX character class.

For instance, [**a-oq-z\cC.**] matches any lowercase latin letter except **p**, as well as control sequences (see below for a description of **\c**).

Quantifiers (repetition).

? 0 or 1, greedy.

?? 0 or 1, lazy.

***** 0 or more, greedy.

***?** 0 or more, lazy.

+ 1 or more, greedy.

+? 1 or more, lazy.

{n} Exactly *n*.

{n,} *n* or more, greedy.

{n,}? *n* or more, lazy.

{n,m} At least *n*, no more than *m*, greedy.

{n,m}? At least *n*, no more than *m*, lazy.

anchors and simple assertions.

\b Word boundary: either the previous token is matched by **\w** and the next by **\W**, or the opposite. For this purpose, the ends of the token list are considered as **\W**.

\B Not a word boundary: between two **\w** tokens or two **\W** tokens (including the boundary).

^ or **\A** Start of the subject token list.

\$, **\Z** or **\z** End of the subject token list.

\G Start of the current match. This is only different from **^** in the case of multiple matches: for instance **\regex_count:nnN { \G a } { aaba } \l_tmpa_int** yields 2, but replacing **\G** by **^** would result in **\l_tmpa_int** holding the value 1.

Alternation and capturing groups.

A|B|C Either one of **A**, **B**, or **C**.

(...) Capturing group.

(?:...) Non-capturing group.

(?*...*) Non-capturing group which resets the group number for capturing groups in each alternative. The following group is numbered with the first unused group number.

The `\c` escape sequence allows to test the category code of tokens, and match control sequences. Each character category is represented by a single uppercase letter:

- C for control sequences;
- B for begin-group tokens;
- E for end-group tokens;
- M for math shift;
- T for alignment tab tokens;
- P for macro parameter tokens;
- U for superscript tokens (up);
- D for subscript tokens (down);
- S for spaces;
- L for letters;
- O for others; and
- A for active characters.

The `\c` escape sequence is used as follows.

`\c{<regex>}` A control sequence whose *cname* matches the *<regex>*, anchored at the beginning and end, so that `\c{begin}` matches exactly `\begin`, and nothing else.

`\cX` Applies to the next object, which can be a character, character property, class, or group, and forces this object to only match tokens with category **X** (any of CBEMTPUDSLOA). For instance, `\cL[A-Z\d]` matches uppercase letters and digits of category code letter, `\cC.` matches any control sequence, and `\cO(abc)` matches `abc` where each character has category other.

`\c[XYZ]` Applies to the next object, and forces it to only match tokens with category **X**, **Y**, or **Z** (each being any of CBEMTPUDSLOA). For instance, `\c[LSO](..)` matches two tokens of category letter, space, or other.

`\c[^XYZ]` Applies to the next object and prevents it from matching any token with category **X**, **Y**, or **Z** (each being any of CBEMTPUDSLOA). For instance, `\c[^O]\d` matches digits which have any category different from other.

The category code tests can be used inside classes; for instance, `[\cO\d \c[LO][A-F]]` matches what \TeX considers as hexadecimal digits, namely digits with category other, or uppercase letters from **A** to **F** with category either letter or other. Within a group affected by a category code test, the outer test can be overridden by a nested test: for instance, `\cL(ab\cO*cd)` matches `ab*cd` where all characters are of category letter, except `*` which has category other.

The `\u` escape sequence allows to insert the contents of a token list directly into a regular expression or a replacement, avoiding the need to escape special characters.

Namely, `\u{<tl var name>}` matches the exact contents of the token list `<tl var>`. Within a `\c{...}` control sequence matching, the `\u` escape sequence only expands its argument once, in effect performing `\tl_to_str:v`. Quantifiers are not supported directly: use a group.

The option `(?i)` makes the match case insensitive (identifying A–Z with a–z; no Unicode support yet). This applies until the end of the group in which it appears, and can be reverted using `(?-i)`. For instance, in `(?i)(a(?-i)b|c)d`, the letters `a` and `d` are affected by the `i` option. Characters within ranges and classes are affected individually: `(?i)[Y-\]` is equivalent to `[YZ\[\]yz]`, and `(?i)[^aeiou]` matches any character which is not a vowel. Neither character properties, nor `\c{...}` nor `\u{...}` are affected by the `i` option.

In character classes, only `[`, `^`, `-`, `]`, `\` and spaces are special, and should be escaped. Other non-alphanumeric characters can still be escaped without harm. Any escape sequence which matches a single character (`\d`, `\D`, *etc.*) is supported in character classes. If the first character is `^`, then the meaning of the character class is inverted; `^` appearing anywhere else in the range is not special. If the first character (possibly following a leading `^`) is `]` then it does not need to be escaped since ending the range there would make it empty. Ranges of characters can be expressed using `-`, for instance, `[\D 0–5]` and `[^6–9]` are equivalent.

Capturing groups are a means of extracting information about the match. Parenthesized groups are labelled in the order of their opening parenthesis, starting at 1. The contents of those groups corresponding to the “best” match (leftmost longest) can be extracted and stored in a sequence of token lists using for instance `\regex_extract_once:nnTF`.

The `\K` escape sequence resets the beginning of the match to the current position in the token list. This only affects what is reported as the full match. For instance,

```
\regex_extract_all:nnN { a \K . } { a123aaxyz } \l_foo_seq
```

results in `\l_foo_seq` containing the items `{1}` and `{a}`: the true matches are `{a1}` and `{aa}`, but they are trimmed by the use of `\K`. The `\K` command does not affect capturing groups: for instance,

```
\regex_extract_once:nnN { (. \K c)+ \d } { acbc3 } \l_foo_seq
```

results in `\l_foo_seq` containing the items `{c3}` and `{bc}`: the true match is `{acbc3}`, with first submatch `{bc}`, but `\K` resets the beginning of the match to the last position where it appears.

2 Syntax of the replacement text

Most of the features described in regular expressions do not make sense within the replacement text. Backslash introduces various special constructions, described further below:

- `\0` is the whole match;
- `\1` is the submatch that was matched by the first (capturing) group `(...)`; similarly for `\2`, ..., `\9` and `\g{<number>}`;
- `_` inserts a space (spaces are ignored when not escaped);

- `\a, \e, \f, \n, \r, \t, \xhh, \x{hhh}` correspond to single characters as in regular expressions;
- `\c{<cs name>}` inserts a control sequence;
- `\c{<category>}<character>` (see below);
- `\u{<tl var name>}` inserts the contents of the `<tl var>` (see below).

Characters other than backslash and space are simply inserted in the result (but since the replacement text is first converted to a string, one should also escape characters that are special for T_EX, for instance use `\#`). Non-alphanumeric characters can always be safely escaped with a backslash.

For instance,

```
\tl_set:Nn \l_my_tl { Hello,~world! }
\regex_replace_all:nnN { ([er]?l|o) . } { (\0--\1) } \l_my_tl
```

results in `\l_my_tl` holding `H(e1l--e1)(o,--o) w(or--o)(ld--l)!`

The submatches are numbered according to the order in which the opening parenthesis of capturing groups appear in the regular expression to match. The n -th submatch is empty if there are fewer than n capturing groups or for capturing groups that appear in alternatives that were not used for the match. In case a capturing group matches several times during a match (due to quantifiers) only the last match is used in the replacement text. Submatches always keep the same category codes as in the original token list.

The characters inserted by the replacement have category code 12 (other) by default, with the exception of space characters. Spaces inserted through `_` have category code 10, while spaces inserted through `\x20` or `\x{20}` have category code 12. The escape sequence `\c` allows to insert characters with arbitrary category codes, as well as control sequences.

`\cX(...)` Produces the characters “...” with category `X`, which must be one of `CBEMTPUDSLOA` as in regular expressions. Parentheses are optional for a single character (which can be an escape sequence). When nested, the innermost category code applies, for instance `\cL(Hello\cS\ world)!` gives this text with standard category codes.

`\c{<text>}` Produces the control sequence with csname `<text>`. The `<text>` may contain references to the submatches `\0`, `\1`, and so on, as in the example for `\u` below.

The escape sequence `\u{<tl var name>}` allows to insert the contents of the token list with name `<tl var name>` directly into the replacement, giving an easier control of category codes. When nested in `\c{...}` and `\u{...}` constructions, the `\u` and `\c` escape sequences perform `\tl_to_str:v`, namely extract the value of the control sequence and turn it into a string. Matches can also be used within the arguments of `\c` and `\u`. For instance,

```
\tl_set:Nn \l_my_one_tl { first }
\tl_set:Nn \l_my_two_tl { \emph{second} }
\tl_set:Nn \l_my_tl { one , two , one , one }
\regex_replace_all:nnN { [,]+ } { \u{1_my_\0_tl} } \l_my_tl
```

results in `\l_my_tl` holding `first,\emph{second},first,first`.

3 Pre-compiling regular expressions

If a regular expression is to be used several times, it is better to compile it once rather than doing it each time the regular expression is used. The compiled regular expression is stored in a variable. All of the `l3regex` module's functions can be given their regular expression argument either as an explicit string or as a compiled regular expression.

`\regex_new:N`

New: 2017-05-26

`\regex_new:N` $\langle regex\ var \rangle$

Creates a new $\langle regex\ var \rangle$ or raises an error if the name is already taken. The declaration is global. The $\langle regex\ var \rangle$ is initially such that it never matches.

`\regex_set:Nn`
`\regex_gset:Nn`
`\regex_const:Nn`

New: 2017-05-26

`\regex_set:Nn` $\langle regex\ var \rangle$ $\{ \langle regex \rangle \}$

Stores a compiled version of the $\langle regular\ expression \rangle$ in the $\langle regex\ var \rangle$. For instance, this function can be used as

```
\regex_new:N \l_my_regex
\regex_set:Nn \l_my_regex { my\ (simple\ )? reg(ex|ular\ expression) }
```

The assignment is local for `\regex_set:Nn` and global for `\regex_gset:Nn`. Use `\regex_const:Nn` for compiled expressions which never change.

`\regex_show:n`
`\regex_show:N`

New: 2017-05-26

`\regex_show:n` $\{ \langle regex \rangle \}$

Shows how `l3regex` interprets the $\langle regex \rangle$. For instance, `\regex_show:n {\A X|Y}` shows

```
+--branch
  anchor at start (\A)
  char code 88
+--branch
  char code 89
```

indicating that the anchor `\A` only applies to the first branch: the second branch is not anchored to the beginning of the match.

4 Matching

All regular expression functions are available in both `:n` and `:N` variants. The former require a “standard” regular expression, while the later require a compiled expression as generated by `\regex_(g)set:Nn`.

`\regex_match:nnTF`
`\regex_match:NnTF`

New: 2017-05-26

`\regex_match:nnTF` $\{ \langle regex \rangle \}$ $\{ \langle token\ list \rangle \}$ $\{ \langle true\ code \rangle \}$ $\{ \langle false\ code \rangle \}$

Tests whether the $\langle regular\ expression \rangle$ matches any part of the $\langle token\ list \rangle$. For instance,

```
\regex_match:nnTF { b [cde]* } { abedcdx } { TRUE } { FALSE }
\regex_match:nnTF { [b-dq-w] } { example } { TRUE } { FALSE }
```

leaves TRUE then FALSE in the input stream.

<code>\regex_count:nnN</code>
<code>\regex_count:NnN</code>
New: 2017-05-26

`\regex_count:nnN` $\{ \langle \textit{regex} \rangle \}$ $\{ \langle \textit{token list} \rangle \}$ $\langle \textit{int var} \rangle$

Sets $\langle \textit{int var} \rangle$ within the current TeX group level equal to the number of times $\langle \textit{regular expression} \rangle$ appears in $\langle \textit{token list} \rangle$. The search starts by finding the left-most longest match, respecting greedy and lazy (non-greedy) operators. Then the search starts again from the character following the last character of the previous match, until reaching the end of the token list. Infinite loops are prevented in the case where the regular expression can match an empty token list: then we count one match between each pair of characters. For instance,

```
\int_new:N \l_foo_int
\regex_count:nnN { (b+|c) } { abbababcb } \l_foo_int
```

results in `\l_foo_int` taking the value 5.

5 Submatch extraction

<code>\regex_extract_once:nnN</code>
<code>\regex_extract_once:nnNTF</code>
<code>\regex_extract_once:NnN</code>
<code>\regex_extract_once:NnNTF</code>
New: 2017-05-26

`\regex_extract_once:nnN` $\{ \langle \textit{regex} \rangle \}$ $\{ \langle \textit{token list} \rangle \}$ $\langle \textit{seq var} \rangle$
`\regex_extract_once:nnNTF` $\{ \langle \textit{regex} \rangle \}$ $\{ \langle \textit{token list} \rangle \}$ $\langle \textit{seq var} \rangle$ $\{ \langle \textit{true code} \rangle \}$ $\{ \langle \textit{false code} \rangle \}$

Finds the first match of the $\langle \textit{regular expression} \rangle$ in the $\langle \textit{token list} \rangle$. If it exists, the match is stored as the first item of the $\langle \textit{seq var} \rangle$, and further items are the contents of capturing groups, in the order of their opening parenthesis. The $\langle \textit{seq var} \rangle$ is assigned locally. If there is no match, the $\langle \textit{seq var} \rangle$ is cleared. The testing versions insert the $\langle \textit{true code} \rangle$ into the input stream if a match was found, and the $\langle \textit{false code} \rangle$ otherwise.

For instance, assume that you type

```
\regex_extract_once:nnNTF { \A(La)?TeX(!*)\Z } { LaTeX!!! } \l_foo_seq
{ true } { false }
```

Then the regular expression (anchored at the start with `\A` and at the end with `\Z`) must match the whole token list. The first capturing group, $(La)?$, matches `La`, and the second capturing group, $(!*)$, matches `!!!`. Thus, `\l_foo_seq` contains as a result the items `{LaTeX!!!}`, `{La}`, and `{!!!}`, and the `true` branch is left in the input stream. Note that the n -th item of `\l_foo_seq`, as obtained using `\seq_item:Nn`, correspond to the submatch numbered $(n - 1)$ in functions such as `\regex_replace_once:nnN`.

<code>\regex_extract_all:nnN</code>
<code>\regex_extract_all:nnNTF</code>
<code>\regex_extract_all:NnN</code>
<code>\regex_extract_all:NnNTF</code>
New: 2017-05-26

`\regex_extract_all:nnN` $\{ \langle \textit{regex} \rangle \}$ $\{ \langle \textit{token list} \rangle \}$ $\langle \textit{seq var} \rangle$
`\regex_extract_all:nnNTF` $\{ \langle \textit{regex} \rangle \}$ $\{ \langle \textit{token list} \rangle \}$ $\langle \textit{seq var} \rangle$ $\{ \langle \textit{true code} \rangle \}$ $\{ \langle \textit{false code} \rangle \}$

Finds all matches of the $\langle \textit{regular expression} \rangle$ in the $\langle \textit{token list} \rangle$, and stores all the submatch information in a single sequence (concatenating the results of multiple `\regex_extract_once:nnN` calls). The $\langle \textit{seq var} \rangle$ is assigned locally. If there is no match, the $\langle \textit{seq var} \rangle$ is cleared. The testing versions insert the $\langle \textit{true code} \rangle$ into the input stream if a match was found, and the $\langle \textit{false code} \rangle$ otherwise. For instance, assume that you type

```
\regex_extract_all:nnNTF { \w+ } { Hello,~world! } \l_foo_seq
{ true } { false }
```

Then the regular expression matches twice, the resulting sequence contains the two items `{Hello}` and `{world}`, and the `true` branch is left in the input stream.

```
\regex_split:nnN
\regex_split:nnNTF
\regex_split:NnN
\regex_split:NnNTF
```

New: 2017-05-26

```
\regex_split:nnN {<regular expression>} {<token list>} <seq var>
\regex_split:nnNTF {<regular expression>} {<token list>} <seq var> {<true code>}
{<false code>}
```

Splits the *<token list>* into a sequence of parts, delimited by matches of the *<regular expression>*. If the *<regular expression>* has capturing groups, then the token lists that they match are stored as items of the sequence as well. The assignment to *<seq var>* is local. If no match is found the resulting *<seq var>* has the *<token list>* as its sole item. If the *<regular expression>* matches the empty token list, then the *<token list>* is split into single tokens. The testing versions insert the *<true code>* into the input stream if a match was found, and the *<false code>* otherwise. For example, after

```
\seq_new:N \l_path_seq
\regex_split:nnNTF { / } { the/path/for/this/file.tex } \l_path_seq
{ true } { false }
```

the sequence `\l_path_seq` contains the items `{the}`, `{path}`, `{for}`, `{this}`, and `{file.tex}`, and the `true` branch is left in the input stream.

6 Replacement

```
\regex_replace_once:nnN
\regex_replace_once:nnNTF
\regex_replace_once:NnN
\regex_replace_once:NnNTF
```

New: 2017-05-26

```
\regex_replace_once:nnN {<regular expression>} {<replacement>} <tl var>
\regex_replace_once:nnNTF {<regular expression>} {<replacement>} <tl var> {<true
code>} {<false code>}
```

Searches for the *<regular expression>* in the *<token list>* and replaces the first match with the *<replacement>*. The result is assigned locally to *<tl var>*. In the *<replacement>*, `\0` represents the full match, `\1` represent the contents of the first capturing group, `\2` of the second, *etc.*

```
\regex_replace_all:nnN
\regex_replace_all:nnNTF
\regex_replace_all:NnN
\regex_replace_all:NnNTF
```

New: 2017-05-26

```
\regex_replace_all:nnN {<regular expression>} {<replacement>} <tl var>
\regex_replace_all:nnNTF {<regular expression>} {<replacement>} <tl var> {<true
code>} {<false code>}
```

Replaces all occurrences of the *<regular expression>* in the *<token list>* by the *<replacement>*, where `\0` represents the full match, `\1` represent the contents of the first capturing group, `\2` of the second, *etc.* Every match is treated independently, and matches cannot overlap. The result is assigned locally to *<tl var>*.

7 Constants and variables

```
\l_tmpa_regex
\l_tmpb_regex
```

New: 2017-12-11

Scratch regex for local assignment. These are never used by the kernel code, and so are safe for use with any L^AT_EX3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.

```
\g_tmpa_regex
\g_tmpb_regex
```

New: 2017-12-11

Scratch regex for global assignment. These are never used by the kernel code, and so are safe for use with any L^AT_EX3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.

8 Bugs, misfeatures, future work, and other possibilities

The following need to be done now.

- Rewrite the documentation in a more ordered way, perhaps add a BNF?
Additional error-checking to come.
- Clean up the use of messages.
- Cleaner error reporting in the replacement phase.
- Add tracing information.
- Detect attempts to use back-references and other non-implemented syntax.
- Test for the maximum register `\c_max_register_int`.
- Find out whether the fact that `\W` and friends match the end-marker leads to bugs.
Possibly update `__regex_item_reverse:n`.
- The empty `cs` should be matched by `\c{}`, not by `\c{csname.?endcsname\s?}`.

Code improvements to come.

- Shift arrays so that the useful information starts at position 1.
- Only build `\c{...}` once.
- Use arrays for the left and right state stacks when compiling a regex.
- Should `__regex_action_free_group:n` only be used for greedy `{n,}` quantifier? (I think not.)
- Quantifiers for `\u` and assertions.
- When matching, keep track of an explicit stack of `curr_state` and `curr_submatches`.
- If possible, when a state is reused by the same thread, kill other subthreads.
- Use an array rather than `\l__regex_balance_tl` to build the function `__regex_replacement_balance_one_match:n`.
- Reduce the number of epsilon-transitions in alternatives.
- Optimize simple strings: use less states (`abcade` should give two states, for `abc` and `ade`). [Does that really make sense?]
- Optimize groups with no alternative.
- Optimize states with a single `__regex_action_free:n`.
- Optimize the use of `__regex_action_success:` by inserting it in state 2 directly instead of having an extra transition.
- Optimize the use of `\int_step_...` functions.

- Groups don't capture within regexes for csnames; optimize and document.
- Better “show” for anchors, properties, and catcode tests.
- Does \K really need a new state for itself?
- When compiling, use a boolean `in_cs` and less magic numbers.
- Instead of checking whether the character is special or alphanumeric using its character code, check if it is special in regexes with `\cs_if_exist` tests.

The following features are likely to be implemented at some point in the future.

- General look-ahead/behind assertions.
- Regex matching on external files.
- Conditional subpatterns with look ahead/behind: “if what follows is [...], then [...]”.
- `(*..)` and `(?..)` sequences to set some options.
- UTF-8 mode for pdfTeX.
- Newline conventions are not done. In particular, we should have an option for `.` not to match newlines. Also, `\A` should differ from `^`, and `\Z`, `\z` and `$` should differ.
- Unicode properties: `\p{..}` and `\P{..}`; `\X` which should match any “extended” Unicode sequence. This requires to manipulate a lot of data, probably using tree-boxes.
- Provide a syntax such as `\ur{1_my_regex}` to use an already-compiled regex in a more complicated regex. This makes regexes more easily composable.
- Allowing `\u{1_my_tl}` in more places, for instance as the number of repetitions in a quantifier.

The following features of PCRE or Perl may or may not be implemented.

- Callout with `(?C...)` or other syntax: some internal code changes make that possible, and it can be useful for instance in the replacement code to stop a regex replacement when some marker has been found; this raises the question of a potential `\regex_break`: and then of playing well with `\tl_map_break`: called from within the code in a regex. It also raises the question of nested calls to the regex machinery, which is a problem since `\fontdimen` are global.
- Conditional subpatterns (other than with a look-ahead or look-behind condition): this is non-regular, isn't it?
- Named subpatterns: TeX programmers have lived so far without any need for named macro parameters.

The following features of PCRE or Perl will definitely not be implemented.

- Back-references: non-regular feature, this requires backtracking, which is prohibitively slow.

- Recursion: this is a non-regular feature.
- Atomic grouping, possessive quantifiers: those tools, mostly meant to fix catastrophic backtracking, are unnecessary in a non-backtracking algorithm, and difficult to implement.
- Subroutine calls: this syntactic sugar is difficult to include in a non-backtracking algorithm, in particular because the corresponding group should be treated as atomic.
- Backtracking control verbs: intrinsically tied to backtracking.
- `\ddd`, matching the character with octal code `ddd`: we already have `\x{...}` and the syntax is confusingly close to what we could have used for backreferences (`\1`, `\2`, ...), making it harder to produce useful error message.
- `\cx`, similar to $\mathrm{T}_{\mathrm{E}}\mathrm{X}$'s own `\^^x`.
- Comments: $\mathrm{T}_{\mathrm{E}}\mathrm{X}$ already has its own system for comments.
- `\Q...\E` escaping: this would require to read the argument verbatim, which is not in the scope of this module.
- `\C` single byte in UTF-8 mode: $\mathrm{X}_{\mathrm{Y}}\mathrm{T}_{\mathrm{E}}\mathrm{X}$ and $\mathrm{LuaT}_{\mathrm{E}}\mathrm{X}$ serve us characters directly, and splitting those into bytes is tricky, encoding dependent, and most likely not useful anyways.

Part XXIX

The l3box package

Boxes

There are three kinds of box operations: horizontal mode denoted with prefix `\hbox_`, vertical mode with prefix `\vbox_`, and the generic operations working in both modes with prefix `\box_`.

1 Creating and initialising boxes

<code>\box_new:N</code>	<code>\box_new:N <box></code>
<code>\box_new:c</code>	Creates a new <code><box></code> or raises an error if the name is already taken. The declaration is global. The <code><box></code> is initially void.

<code>\box_clear:N</code>	<code>\box_clear:N <box></code>
<code>\box_clear:c</code>	Clears the content of the <code><box></code> by setting the box equal to <code>\c_empty_box</code> .
<code>\box_gclear:N</code>	
<code>\box_gclear:c</code>	

<code>\box_clear_new:N</code>	<code>\box_clear_new:N <box></code>
<code>\box_clear_new:c</code>	Ensures that the <code><box></code> exists globally by applying <code>\box_new:N</code> if necessary, then applies <code>\box_(g)clear:N</code> to leave the <code><box></code> empty.
<code>\box_gclear_new:N</code>	
<code>\box_gclear_new:c</code>	

<code>\box_set_eq:NN</code>	<code>\box_set_eq:NN <box₁> <box₂></code>
<code>\box_set_eq:(cN Nc cc)</code>	Sets the content of <code><box₁></code> equal to that of <code><box₂></code> .
<code>\box_gset_eq:NN</code>	
<code>\box_gset_eq:(cN Nc cc)</code>	

<code>\box_if_exist_p:N *</code>	<code>\box_if_exist_p:N <box></code>
<code>\box_if_exist_p:c *</code>	<code>\box_if_exist:NTF <box> {(true code)} {(false code)}</code>
<code>\box_if_exist:NTF *</code>	Tests whether the <code><box></code> is currently defined. This does not check that the <code><box></code> really is a box.
<code>\box_if_exist:cTF *</code>	

New: 2012-03-03

2 Using boxes

<code>\box_use:N</code>	<code>\box_use:N <box></code>
<code>\box_use:c</code>	Inserts the current content of the <code><box></code> onto the current list for typesetting. An error is raised if the variable does not exist or if it is invalid.

T_EXhackers note: This is the T_EX primitive `\copy`.

<hr/> <code>\box_move_right:nn</code> <hr/>	<code>\box_move_right:nn {<dimexpr>} {<box function>}</code>
<code>\box_move_left:nn</code> <hr/>	This function operates in vertical mode, and inserts the material specified by the <i><box function></i> such that its reference point is displaced horizontally by the given <i><dimexpr></i> from the reference point for typesetting, to the right or left as appropriate. The <i><box function></i> should be a box operation such as <code>\box_use:N \<box></code> or a “raw” box specification such as <code>\vbox:n { xyz }</code> .

<hr/> <code>\box_move_up:nn</code> <hr/>	<code>\box_move_up:nn {<dimexpr>} {<box function>}</code>
<code>\box_move_down:nn</code> <hr/>	This function operates in horizontal mode, and inserts the material specified by the <i><box function></i> such that its reference point is displaced vertically by the given <i><dimexpr></i> from the reference point for typesetting, up or down as appropriate. The <i><box function></i> should be a box operation such as <code>\box_use:N \<box></code> or a “raw” box specification such as <code>\vbox:n { xyz }</code> .

3 Measuring and setting box dimensions

<hr/> <code>\box_dp:N</code> <hr/>	<code>\box_dp:N <box></code>
<code>\box_dp:c</code> <hr/>	Calculates the depth (below the baseline) of the <i><box></i> in a form suitable for use in a <i><dimension expression></i> .

TeXhackers note: This is the TeX primitive `\dp`.

<hr/> <code>\box_ht:N</code> <hr/>	<code>\box_ht:N <box></code>
<code>\box_ht:c</code> <hr/>	Calculates the height (above the baseline) of the <i><box></i> in a form suitable for use in a <i><dimension expression></i> .

TeXhackers note: This is the TeX primitive `\ht`.

<hr/> <code>\box_wd:N</code> <hr/>	<code>\box_wd:N <box></code>
<code>\box_wd:c</code> <hr/>	Calculates the width of the <i><box></i> in a form suitable for use in a <i><dimension expression></i> .

TeXhackers note: This is the TeX primitive `\wd`.

<hr/> <code>\box_set_dp:Nn</code> <hr/>	<code>\box_set_dp:Nn <box> {<dimension expression>}</code>
<code>\box_set_dp:cn</code> <hr/>	Set the depth (below the baseline) of the <i><box></i> to the value of the <i>{<dimension expression>}</i> .
<code>\box_gset_dp:Nn</code> <hr/>	
<code>\box_gset_dp:cn</code> <hr/>	

Updated: 2019-01-22

<hr/> <code>\box_set_ht:Nn</code> <hr/>	<code>\box_set_ht:Nn <box> {<dimension expression>}</code>
<code>\box_set_ht:cn</code> <hr/>	Set the height (above the baseline) of the <i><box></i> to the value of the <i>{<dimension expression>}</i> .
<code>\box_gset_ht:Nn</code> <hr/>	
<code>\box_gset_ht:cn</code> <hr/>	

Updated: 2019-01-22

<code>\box_set_wd:Nn</code>	<code>\box_set_wd:Nn <box> {<dimension expression>}</code>
<code>\box_set_wd:cn</code>	
<code>\box_gset_wd:Nn</code>	Set the width of the <code><box></code> to the value of the <code>{<dimension expression>}</code> .
<code>\box_gset_wd:cn</code>	

Updated: 2019-01-22

4 Box conditionals

<code>\box_if_empty_p:N</code> *	<code>\box_if_empty_p:N <box></code>
<code>\box_if_empty_p:c</code> *	<code>\box_if_empty:NTF <box> {<true code>} {<false code>}</code>
<code>\box_if_empty:NTF</code> *	
<code>\box_if_empty:cTF</code> *	Tests if <code><box></code> is a empty (equal to <code>\c_empty_box</code>).

<code>\box_if_horizontal_p:N</code> *	<code>\box_if_horizontal_p:N <box></code>
<code>\box_if_horizontal_p:c</code> *	<code>\box_if_horizontal:NTF <box> {<true code>} {<false code>}</code>
<code>\box_if_horizontal:NTF</code> *	
<code>\box_if_horizontal:cTF</code> *	Tests if <code><box></code> is a horizontal box.

<code>\box_if_vertical_p:N</code> *	<code>\box_if_vertical_p:N <box></code>
<code>\box_if_vertical_p:c</code> *	<code>\box_if_vertical:NTF <box> {<true code>} {<false code>}</code>
<code>\box_if_vertical:NTF</code> *	
<code>\box_if_vertical:cTF</code> *	Tests if <code><box></code> is a vertical box.

5 The last box inserted

<code>\box_set_to_last:N</code>	<code>\box_set_to_last:N <box></code>
<code>\box_set_to_last:c</code>	
<code>\box_gset_to_last:N</code>	Sets the <code><box></code> equal to the last item (box) added to the current partial list, removing the item from the list at the same time. When applied to the main vertical list, the <code><box></code> is always void as it is not possible to recover the last added item.
<code>\box_gset_to_last:c</code>	

6 Constant boxes

<code>\c_empty_box</code>	This is a permanently empty box, which is neither set as horizontal nor vertical.
---------------------------	---

Updated: 2012-11-04

T_EXhackers note: At the T_EX level this is a void box.

7 Scratch boxes

`\l_tmpa_box`
`\l_tmpb_box`

Updated: 2012-11-04

Scratch boxes for local assignment. These are never used by the kernel code, and so are safe for use with any L^AT_EX3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.

`\g_tmpa_box`
`\g_tmpb_box`

Scratch boxes for global assignment. These are never used by the kernel code, and so are safe for use with any L^AT_EX3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.

8 Viewing box contents

`\box_show:N`
`\box_show:c`

Updated: 2012-05-11

`\box_show:N` $\langle box \rangle$

Shows full details of the content of the $\langle box \rangle$ in the terminal.

`\box_show:Nnn`
`\box_show:cnn`

New: 2012-05-11

`\box_show:Nnn` $\langle box \rangle$ $\{\langle intexpr_1 \rangle\}$ $\{\langle intexpr_2 \rangle\}$

Display the contents of $\langle box \rangle$ in the terminal, showing the first $\langle intexpr_1 \rangle$ items of the box, and descending into $\langle intexpr_2 \rangle$ group levels.

`\box_log:N`
`\box_log:c`

New: 2012-05-11

`\box_log:N` $\langle box \rangle$

Writes full details of the content of the $\langle box \rangle$ to the log.

`\box_log:Nnn`
`\box_log:cnn`

New: 2012-05-11

`\box_log:Nnn` $\langle box \rangle$ $\{\langle intexpr_1 \rangle\}$ $\{\langle intexpr_2 \rangle\}$

Writes the contents of $\langle box \rangle$ to the log, showing the first $\langle intexpr_1 \rangle$ items of the box, and descending into $\langle intexpr_2 \rangle$ group levels.

9 Boxes and color

All L^AT_EX3 boxes are “color safe”: a color set inside the box stops applying after the end of the box has occurred.

10 Horizontal mode boxes

`\hbox:n`

Updated: 2017-04-05

`\hbox:n` $\{\langle contents \rangle\}$

Typesets the $\langle contents \rangle$ into a horizontal box of natural width and then includes this box in the current list for typesetting.

<hr/> <code>\hbox_to_wd:nn</code> <hr/>	<code>\hbox_to_wd:nn {<dimexpr>} {<contents>}</code>
Updated: 2017-04-05 <hr/>	Typesets the $\langle contents \rangle$ into a horizontal box of width $\langle dimexpr \rangle$ and then includes this box in the current list for typesetting.
<hr/> <code>\hbox_to_zero:n</code> <hr/>	<code>\hbox_to_zero:n {<contents>}</code>
Updated: 2017-04-05 <hr/>	Typesets the $\langle contents \rangle$ into a horizontal box of zero width and then includes this box in the current list for typesetting.
<hr/> <code>\hbox_set:Nn</code> <code>\hbox_set:cn</code> <code>\hbox_gset:Nn</code> <code>\hbox_gset:cn</code> <hr/>	<code>\hbox_set:Nn <box> {<contents>}</code> Typesets the $\langle contents \rangle$ at natural width and then stores the result inside the $\langle box \rangle$.
Updated: 2017-04-05 <hr/>	
<hr/> <code>\hbox_set_to_wd:Nnn</code> <code>\hbox_set_to_wd:cnn</code> <code>\hbox_gset_to_wd:Nnn</code> <code>\hbox_gset_to_wd:cnn</code> <hr/>	<code>\hbox_set_to_wd:Nnn <box> {<dimexpr>} {<contents>}</code> Typesets the $\langle contents \rangle$ to the width given by the $\langle dimexpr \rangle$ and then stores the result inside the $\langle box \rangle$.
Updated: 2017-04-05 <hr/>	
<hr/> <code>\hbox_overlap_center:n</code> <hr/>	<code>\hbox_overlap_center:n {<contents>}</code>
New: 2020-08-25 <hr/>	Typesets the $\langle contents \rangle$ into a horizontal box of zero width such that material protrudes equally to both sides of the insertion point.
<hr/> <code>\hbox_overlap_right:n</code> <hr/>	<code>\hbox_overlap_right:n {<contents>}</code>
Updated: 2017-04-05 <hr/>	Typesets the $\langle contents \rangle$ into a horizontal box of zero width such that material protrudes to the right of the insertion point.
<hr/> <code>\hbox_overlap_left:n</code> <hr/>	<code>\hbox_overlap_left:n {<contents>}</code>
Updated: 2017-04-05 <hr/>	Typesets the $\langle contents \rangle$ into a horizontal box of zero width such that material protrudes to the left of the insertion point.
<hr/> <code>\hbox_set:Nw</code> <code>\hbox_set:cw</code> <code>\hbox_set_end:</code> <code>\hbox_gset:Nw</code> <code>\hbox_gset:cw</code> <code>\hbox_gset_end:</code> <hr/>	<code>\hbox_set:Nw <box> <contents> \hbox_set_end:</code> Typesets the $\langle contents \rangle$ at natural width and then stores the result inside the $\langle box \rangle$. In contrast to <code>\hbox_set:Nn</code> this function does not absorb the argument when finding the $\langle content \rangle$, and so can be used in circumstances where the $\langle content \rangle$ may not be a simple argument.
Updated: 2017-04-05 <hr/>	
<hr/> <code>\hbox_set_to_wd:Nnw</code> <code>\hbox_set_to_wd:cnw</code> <code>\hbox_gset_to_wd:Nnw</code> <code>\hbox_gset_to_wd:cnw</code> <hr/>	<code>\hbox_set_to_wd:Nnw <box> {<dimexpr>} <contents> \hbox_set_end:</code> Typesets the $\langle contents \rangle$ to the width given by the $\langle dimexpr \rangle$ and then stores the result inside the $\langle box \rangle$. In contrast to <code>\hbox_set_to_wd:Nnn</code> this function does not absorb the argument when finding the $\langle content \rangle$, and so can be used in circumstances where the $\langle content \rangle$ may not be a simple argument
New: 2017-06-08 <hr/>	

<hr/> <code>\hbox_unpack:N</code> <hr/>	<code>\hbox_unpack:N <box></code>
<code>\hbox_unpack:c</code> <hr/>	Unpacks the content of the horizontal $\langle box \rangle$, retaining any stretching or shrinking applied when the $\langle box \rangle$ was set.

T_EXhackers note: This is the T_EX primitive `\unhcopy`.

11 Vertical mode boxes

Vertical boxes inherit their baseline from their contents. The standard case is that the baseline of the box is at the same position as that of the last item added to the box. This means that the box has no depth unless the last item added to it had depth. As a result most vertical boxes have a large height value and small or zero depth. The exception are `_top` boxes, where the reference point is that of the first item added. These tend to have a large depth and small height, although the latter is typically non-zero.

<hr/> <code>\vbox:n</code> <hr/>	<code>\vbox:n {\langle contents \rangle}</code>
<code>Updated: 2017-04-05</code> <hr/>	Typesets the $\langle contents \rangle$ into a vertical box of natural height and includes this box in the current list for typesetting.

<hr/> <code>\vbox_top:n</code> <hr/>	<code>\vbox_top:n {\langle contents \rangle}</code>
<code>Updated: 2017-04-05</code> <hr/>	Typesets the $\langle contents \rangle$ into a vertical box of natural height and includes this box in the current list for typesetting. The baseline of the box is equal to that of the <i>first</i> item added to the box.

<hr/> <code>\vbox_to_ht:nn</code> <hr/>	<code>\vbox_to_ht:nn {\langle dimexpr \rangle} {\langle contents \rangle}</code>
<code>Updated: 2017-04-05</code> <hr/>	Typesets the $\langle contents \rangle$ into a vertical box of height $\langle dimexpr \rangle$ and then includes this box in the current list for typesetting.

<hr/> <code>\vbox_to_zero:n</code> <hr/>	<code>\vbox_to_zero:n {\langle contents \rangle}</code>
<code>Updated: 2017-04-05</code> <hr/>	Typesets the $\langle contents \rangle$ into a vertical box of zero height and then includes this box in the current list for typesetting.

<hr/> <code>\vbox_set:Nn</code> <hr/>	<code>\vbox_set:Nn <box> {\langle contents \rangle}</code>
<code>\vbox_set:cn</code> <hr/>	Typesets the $\langle contents \rangle$ at natural height and then stores the result inside the $\langle box \rangle$.
<code>\vbox_gset:Nn</code> <hr/>	
<code>\vbox_gset:cn</code> <hr/>	
<code>Updated: 2017-04-05</code> <hr/>	

<hr/> <code>\vbox_set_top:Nn</code> <hr/>	<code>\vbox_set_top:Nn <box> {\langle contents \rangle}</code>
<code>\vbox_set_top:cn</code> <hr/>	Typesets the $\langle contents \rangle$ at natural height and then stores the result inside the $\langle box \rangle$. The baseline of the box is equal to that of the <i>first</i> item added to the box.
<code>\vbox_gset_top:Nn</code> <hr/>	
<code>\vbox_gset_top:cn</code> <hr/>	
<code>Updated: 2017-04-05</code> <hr/>	

<code>\vbox_set_to_ht:Nnn</code>	<code>\vbox_set_to_ht:Nnn <box> {<dimexpr>} {<contents>}</code>
<code>\vbox_set_to_ht:cnn</code>	
<code>\vbox_gset_to_ht:Nnn</code>	Typesets the $\langle contents \rangle$ to the height given by the $\langle dimexpr \rangle$ and then stores the result inside the $\langle box \rangle$.
<code>\vbox_gset_to_ht:cnn</code>	

Updated: 2017-04-05

<code>\vbox_set:Nw</code>	<code>\vbox_set:Nw <box> <contents> \vbox_set_end:</code>
<code>\vbox_set:cw</code>	
<code>\vbox_set_end:</code>	Typesets the $\langle contents \rangle$ at natural height and then stores the result inside the $\langle box \rangle$. In contrast to <code>\vbox_set:Nn</code> this function does not absorb the argument when finding the $\langle content \rangle$, and so can be used in circumstances where the $\langle content \rangle$ may not be a simple argument.
<code>\vbox_gset:Nw</code>	
<code>\vbox_gset:cw</code>	
<code>\vbox_gset_end:</code>	

Updated: 2017-04-05

<code>\vbox_set_to_ht:Nnw</code>	<code>\vbox_set_to_ht:Nnw <box> {<dimexpr>} <contents> \vbox_set_end:</code>
<code>\vbox_set_to_ht:cnw</code>	
<code>\vbox_gset_to_ht:Nnw</code>	Typesets the $\langle contents \rangle$ to the height given by the $\langle dimexpr \rangle$ and then stores the result inside the $\langle box \rangle$. In contrast to <code>\vbox_set_to_ht:Nnn</code> this function does not absorb the argument when finding the $\langle content \rangle$, and so can be used in circumstances where the $\langle content \rangle$ may not be a simple argument
<code>\vbox_gset_to_ht:cnw</code>	

New: 2017-06-08

<code>\vbox_set_split_to_ht:NNn</code>	<code>\vbox_set_split_to_ht:NNn <box₁> <box₂> {<dimexpr>}</code>
<code>\vbox_set_split_to_ht:(cNn Ncn ccn)</code>	
<code>\vbox_gset_split_to_ht:NNn</code>	
<code>\vbox_gset_split_to_ht:(cNn Ncn ccn)</code>	

Updated: 2018-12-29

Sets $\langle box_1 \rangle$ to contain material to the height given by the $\langle dimexpr \rangle$ by removing content from the top of $\langle box_2 \rangle$ (which must be a vertical box).

<code>\vbox_unpack:N</code>	<code>\vbox_unpack:N <box></code>
<code>\vbox_unpack:c</code>	Unpacks the content of the vertical $\langle box \rangle$, retaining any stretching or shrinking applied when the $\langle box \rangle$ was set.

TeXhackers note: This is the TeX primitive `\unvcopy`.

12 Using boxes efficiently

The functions above for using box contents work in exactly the same way as for any other `expl3` variable. However, for efficiency reasons, it is also useful to have functions which *drop* box contents on use. When a box is dropped, the box becomes empty at the group level *where the box was originally set* rather than necessarily *at the current group level*. For example, with

```
\hbox_set:Nn \l_tmpa_box { A }
\group_begin:
  \hbox_set:Nn \l_tmpa_box { B }
\group_begin:
```

```

\box_use_drop:N \l_tmpa_box
\group_end:
\box_show:N \l_tmpa_box
\group_end:
\box_show:N \l_tmpa_box

```

the first use of `\box_show:N` will show an entirely cleared (void) box, and the second will show the letter A in the box.

These functions should be preferred when the content of the box is no longer required after use. Note that due to the unusual scoping behaviour of `drop` functions they may be applied to both local and global boxes: the latter will naturally be set and thus cleared at a global level.

```

\box_use_drop:N
\box_use_drop:c

```

`\box_use_drop:N <box>`

Inserts the current content of the $\langle box \rangle$ onto the current list for typesetting then drops the box content. An error is raised if the variable does not exist or if it is invalid. This function may be applied to local or global boxes.

TeXhackers note: This is the `\box` primitive.

```

\box_set_eq_drop:NN
\box_set_eq_drop:(cN|Nc|cc)

```

New: 2019-01-17

`\box_set_eq_drop:NN <box1> <box2>`

Sets the content of $\langle box_1 \rangle$ equal to that of $\langle box_2 \rangle$, then drops $\langle box_2 \rangle$.

```

\box_gset_eq_drop:NN
\box_gset_eq_drop:(cN|Nc|cc)

```

New: 2019-01-17

`\box_gset_eq_drop:NN <box1> <box2>`

Sets the content of $\langle box_1 \rangle$ globally equal to that of $\langle box_2 \rangle$, then drops $\langle box_2 \rangle$.

```

\hbox_unpack_drop:N
\hbox_unpack_drop:c

```

New: 2019-01-17

`\hbox_unpack_drop:N <box>`

Unpacks the content of the horizontal $\langle box \rangle$, retaining any stretching or shrinking applied when the $\langle box \rangle$ was set. The original $\langle box \rangle$ is then dropped.

TeXhackers note: This is the TeX primitive `\unhbox`.

```

\vbox_unpack_drop:N
\vbox_unpack_drop:c

```

New: 2019-01-17

`\vbox_unpack_drop:N <box>`

Unpacks the content of the vertical $\langle box \rangle$, retaining any stretching or shrinking applied when the $\langle box \rangle$ was set. The original $\langle box \rangle$ is then dropped.

TeXhackers note: This is the TeX primitive `\unvbox`.

13 Affine transformations

Affine transformations are changes which (informally) preserve straight lines. Simple translations are affine transformations, but are better handled in TeX by doing the translation first, then inserting an unmodified box. On the other hand, rotation and resizing

of boxed material can best be handled by modifying boxes. These transformations are described here.

<code>\box_autosize_to_wd_and_ht:Nnn</code>	<code>\box_autosize_to_wd_and_ht:Nnn <box> {<x-size>} {<y-size>}</code>
<code>\box_autosize_to_wd_and_ht:cnn</code>	
<code>\box_gautosize_to_wd_and_ht:Nnn</code>	
<code>\box_gautosize_to_wd_and_ht:cnn</code>	

New: 2017-04-04

Updated: 2019-01-22

Resizes the $\langle box \rangle$ to fit within the given $\langle x-size \rangle$ (horizontally) and $\langle y-size \rangle$ (vertically); both of the sizes are dimension expressions. The $\langle y-size \rangle$ is the height only: it does not include any depth. The updated $\langle box \rangle$ is an `hbox`, irrespective of the nature of the $\langle box \rangle$ before the resizing is applied. The final size of the $\langle box \rangle$ is the smaller of $\{ \langle x-size \rangle \}$ and $\{ \langle y-size \rangle \}$, *i.e.* the result fits within the dimensions specified. Negative sizes cause the material in the $\langle box \rangle$ to be reversed in direction, but the reference point of the $\langle box \rangle$ is unchanged. Thus a negative $\langle y-size \rangle$ results in the $\langle box \rangle$ having a depth dependent on the height of the original and *vice versa*.

<code>\box_autosize_to_wd_and_ht_plus_dp:Nnn</code>	<code>\box_autosize_to_wd_and_ht_plus_dp:Nnn <box> {<x-size>}</code>
<code>\box_autosize_to_wd_and_ht_plus_dp:cnn</code>	<code>{<y-size>}</code>
<code>\box_gautosize_to_wd_and_ht_plus_dp:Nnn</code>	
<code>\box_gautosize_to_wd_and_ht_plus_dp:cnn</code>	

New: 2017-04-04

Updated: 2019-01-22

Resizes the $\langle box \rangle$ to fit within the given $\langle x-size \rangle$ (horizontally) and $\langle y-size \rangle$ (vertically); both of the sizes are dimension expressions. The $\langle y-size \rangle$ is the total vertical size (height plus depth). The updated $\langle box \rangle$ is an `hbox`, irrespective of the nature of the $\langle box \rangle$ before the resizing is applied. The final size of the $\langle box \rangle$ is the smaller of $\{ \langle x-size \rangle \}$ and $\{ \langle y-size \rangle \}$, *i.e.* the result fits within the dimensions specified. Negative sizes cause the material in the $\langle box \rangle$ to be reversed in direction, but the reference point of the $\langle box \rangle$ is unchanged. Thus a negative $\langle y-size \rangle$ results in the $\langle box \rangle$ having a depth dependent on the height of the original and *vice versa*.

<code>\box_resize_to_ht:Nn</code>	<code>\box_resize_to_ht:Nn <box> {<y-size>}</code>
<code>\box_resize_to_ht:cn</code>	
<code>\box_gresize_to_ht:Nn</code>	
<code>\box_gresize_to_ht:cn</code>	

Updated: 2019-01-22

Resizes the $\langle box \rangle$ to $\langle y-size \rangle$ (vertically), scaling the horizontal size by the same amount; $\langle y-size \rangle$ is a dimension expression. The $\langle y-size \rangle$ is the height only: it does not include any depth. The updated $\langle box \rangle$ is an `hbox`, irrespective of the nature of the $\langle box \rangle$ before the resizing is applied. A negative $\langle y-size \rangle$ causes the material in the $\langle box \rangle$ to be reversed in direction, but the reference point of the $\langle box \rangle$ is unchanged. Thus a negative $\langle y-size \rangle$ results in the $\langle box \rangle$ having a depth dependent on the height of the original and *vice versa*.

<code>\box_resize_to_ht_plus_dp:Nn</code>	<code>\box_resize_to_ht_plus_dp:Nn <box> {<y-size>}</code>
<code>\box_resize_to_ht_plus_dp:cn</code>	
<code>\box_gresize_to_ht_plus_dp:Nn</code>	
<code>\box_gresize_to_ht_plus_dp:cn</code>	

Updated: 2019-01-22

Resizes the $\langle box \rangle$ to $\langle y-size \rangle$ (vertically), scaling the horizontal size by the same amount; $\langle y-size \rangle$ is a dimension expression. The $\langle y-size \rangle$ is the total vertical size (height plus depth). The updated $\langle box \rangle$ is an `hbox`, irrespective of the nature of the $\langle box \rangle$ before the resizing is applied. A negative $\langle y-size \rangle$ causes the material in the $\langle box \rangle$ to be reversed in direction, but the reference point of the $\langle box \rangle$ is unchanged. Thus a negative $\langle y-size \rangle$ results in the $\langle box \rangle$ having a depth dependent on the height of the original and *vice versa*.

<code>\box_resize_to_wd:Nn</code>	<code>\box_resize_to_wd:Nn <box> {<x-size>}</code>
<code>\box_resize_to_wd:cn</code>	
<code>\box_gresize_to_wd:Nn</code>	
<code>\box_gresize_to_wd:cn</code>	

Updated: 2019-01-22

Resizes the $\langle box \rangle$ to $\langle x-size \rangle$ (horizontally), scaling the vertical size by the same amount; $\langle x-size \rangle$ is a dimension expression. The updated $\langle box \rangle$ is an `hbox`, irrespective of the nature of the $\langle box \rangle$ before the resizing is applied. A negative $\langle x-size \rangle$ causes the material in the $\langle box \rangle$ to be reversed in direction, but the reference point of the $\langle box \rangle$ is unchanged. Thus a negative $\langle x-size \rangle$ results in the $\langle box \rangle$ having a depth dependent on the height of the original and *vice versa*.

<code>\box_resize_to_wd_and_ht:Nnn</code>	<code>\box_resize_to_wd_and_ht:Nnn <box> {<x-size>} {<y-size>}</code>
<code>\box_resize_to_wd_and_ht:cnn</code>	
<code>\box_gresize_to_wd_and_ht:Nnn</code>	
<code>\box_gresize_to_wd_and_ht:cnn</code>	

New: 2014-07-03

Updated: 2019-01-22

Resizes the $\langle box \rangle$ to $\langle x-size \rangle$ (horizontally) and $\langle y-size \rangle$ (vertically): both of the sizes are dimension expressions. The $\langle y-size \rangle$ is the height only and does not include any depth. The updated $\langle box \rangle$ is an `hbox`, irrespective of the nature of the $\langle box \rangle$ before the resizing is applied. Negative sizes cause the material in the $\langle box \rangle$ to be reversed in direction, but the reference point of the $\langle box \rangle$ is unchanged. Thus a negative $\langle y-size \rangle$ results in the $\langle box \rangle$ having a depth dependent on the height of the original and *vice versa*.

<code>\box_resize_to_wd_and_ht_plus_dp:Nnn</code>	<code>\box_resize_to_wd_and_ht_plus_dp:Nnn <box> {<x-size>} {<y-size>}</code>
<code>\box_resize_to_wd_and_ht_plus_dp:cnn</code>	
<code>\box_gresize_to_wd_and_ht_plus_dp:Nnn</code>	
<code>\box_gresize_to_wd_and_ht_plus_dp:cnn</code>	

New: 2017-04-06

Updated: 2019-01-22

Resizes the $\langle box \rangle$ to $\langle x-size \rangle$ (horizontally) and $\langle y-size \rangle$ (vertically): both of the sizes are dimension expressions. The $\langle y-size \rangle$ is the total vertical size (height plus depth). The updated $\langle box \rangle$ is an `hbox`, irrespective of the nature of the $\langle box \rangle$ before the resizing is applied. Negative sizes cause the material in the $\langle box \rangle$ to be reversed in direction, but the reference point of the $\langle box \rangle$ is unchanged. Thus a negative $\langle y-size \rangle$ results in the $\langle box \rangle$ having a depth dependent on the height of the original and *vice versa*.

<code>\box_rotate:Nn</code>	<code>\box_rotate:Nn <box> {<angle>}</code>
<code>\box_rotate:cn</code>	
<code>\box_grotate:Nn</code>	Rotates the $\langle box \rangle$ by $\langle angle \rangle$ (in degrees) anti-clockwise about its reference point. The reference point of the updated box is moved horizontally such that it is at the left side of the smallest rectangle enclosing the rotated material. The updated $\langle box \rangle$ is an <code>hbox</code> , irrespective of the nature of the $\langle box \rangle$ before the rotation is applied.
<code>\box_grotate:cn</code>	
<hr/> Updated: 2019-01-22 <hr/>	

<code>\box_scale:Nnn</code>	<code>\box_scale:Nnn <box> {<x-scale>} {<y-scale>}</code>
<code>\box_scale:cnn</code>	
<code>\box_gscale:Nnn</code>	Scales the $\langle box \rangle$ by factors $\langle x-scale \rangle$ and $\langle y-scale \rangle$ in the horizontal and vertical directions, respectively (both scales are integer expressions). The updated $\langle box \rangle$ is an <code>hbox</code> , irrespective of the nature of the $\langle box \rangle$ before the scaling is applied. Negative scalings cause the material in the $\langle box \rangle$ to be reversed in direction, but the reference point of the $\langle box \rangle$ is unchanged. Thus a negative $\langle y-scale \rangle$ results in the $\langle box \rangle$ having a depth dependent on the height of the original and <i>vice versa</i> .
<code>\box_gscale:cnn</code>	
<hr/> Updated: 2019-01-22 <hr/>	

14 Primitive box conditionals

<code>\if_hbox:N *</code>	<code>\if_hbox:N <box></code> <code> <true code></code> <code>\else:</code> <code> <false code></code> <code>\fi:</code> Tests is $\langle box \rangle$ is a horizontal box. TeXhackers note: This is the TeX primitive <code>\ifhbox</code> .
---------------------------	---

<code>\if_vbox:N *</code>	<code>\if_vbox:N <box></code> <code> <true code></code> <code>\else:</code> <code> <false code></code> <code>\fi:</code> Tests is $\langle box \rangle$ is a vertical box. TeXhackers note: This is the TeX primitive <code>\ifvbox</code> .
---------------------------	---

<code>\if_box_empty:N *</code>	<code>\if_box_empty:N <box></code> <code> <true code></code> <code>\else:</code> <code> <false code></code> <code>\fi:</code> Tests is $\langle box \rangle$ is an empty (void) box. TeXhackers note: This is the TeX primitive <code>\ifvoid</code> .
--------------------------------	---

Part XXX

The l3coffins package

Coffin code layer

The material in this module provides the low-level support system for coffins. For details about the design concept of a coffin, see the xcoffins module (in the l3experimental bundle).

1 Creating and initialising coffins

`\coffin_new:N`
`\coffin_new:c`

New: 2011-08-17

`\coffin_new:N` $\langle coffin \rangle$

Creates a new $\langle coffin \rangle$ or raises an error if the name is already taken. The declaration is global. The $\langle coffin \rangle$ is initially empty.

`\coffin_clear:N`
`\coffin_clear:c`
`\coffin_gclear:N`
`\coffin_gclear:c`

New: 2011-08-17
Updated: 2019-01-21

`\coffin_clear:N` $\langle coffin \rangle$

Clears the content of the $\langle coffin \rangle$.

`\coffin_set_eq:NN`
`\coffin_set_eq:(Nc|cN|cc)`
`\coffin_gset_eq:NN`
`\coffin_gset_eq:(Nc|cN|cc)`

New: 2011-08-17
Updated: 2019-01-21

`\coffin_set_eq:NN` $\langle coffin_1 \rangle$ $\langle coffin_2 \rangle$

Sets both the content and poles of $\langle coffin_1 \rangle$ equal to those of $\langle coffin_2 \rangle$.

`\coffin_if_exist_p:N` \star
`\coffin_if_exist_p:c` \star
`\coffin_if_exist:N \overline{TF}` \star
`\coffin_if_exist:c \overline{TF}` \star

New: 2012-06-20

`\coffin_if_exist_p:N` $\langle box \rangle$

`\coffin_if_exist:NTF` $\langle box \rangle$ $\{\langle true\ code \rangle\}$ $\{\langle false\ code \rangle\}$

Tests whether the $\langle coffin \rangle$ is currently defined.

2 Setting coffin content and poles

`\hcoffin_set:Nn`
`\hcoffin_set:cn`
`\hcoffin_gset:Nn`
`\hcoffin_gset:cn`

New: 2011-08-17
Updated: 2019-01-21

`\hcoffin_set:Nn` $\langle coffin \rangle$ $\{\langle material \rangle\}$

Typesets the $\langle material \rangle$ in horizontal mode, storing the result in the $\langle coffin \rangle$. The standard poles for the $\langle coffin \rangle$ are then set up based on the size of the typeset material.

```

\hcoffin_set:Nw
\hcoffin_set:cw
\hcoffin_set_end:
\hcoffin_gset:Nw
\hcoffin_gset:cw
\hcoffin_gset_end:

```

New: 2011-09-10
Updated: 2019-01-21

```

\vcoffin_set:Nnn
\vcoffin_set:cnn
\vcoffin_gset:Nnn
\vcoffin_gset:cnn

```

New: 2011-08-17
Updated: 2019-01-21

```

\vcoffin_set:Nnw
\vcoffin_set:cnw
\vcoffin_set_end:
\vcoffin_gset:Nnw
\vcoffin_gset:cnw
\vcoffin_gset_end:

```

New: 2011-09-10
Updated: 2019-01-21

`\hcoffin_set:Nw <coffin> <material> \hcoffin_set_end:`

Typesets the $\langle material \rangle$ in horizontal mode, storing the result in the $\langle coffin \rangle$. The standard poles for the $\langle coffin \rangle$ are then set up based on the size of the typeset material. These functions are useful for setting the entire contents of an environment in a coffin.

`\vcoffin_set:Nnn <coffin> {\<width>} {\<material>}`

Typesets the $\langle material \rangle$ in vertical mode constrained to the given $\langle width \rangle$ and stores the result in the $\langle coffin \rangle$. The standard poles for the $\langle coffin \rangle$ are then set up based on the size of the typeset material.

`\vcoffin_set:Nnw <coffin> {\<width>} <material> \vcoffin_set_end:`

Typesets the $\langle material \rangle$ in vertical mode constrained to the given $\langle width \rangle$ and stores the result in the $\langle coffin \rangle$. The standard poles for the $\langle coffin \rangle$ are then set up based on the size of the typeset material. These functions are useful for setting the entire contents of an environment in a coffin.

```

\coffin_set_horizontal_pole:Nnn
\coffin_set_horizontal_pole:cnn
\coffin_gset_horizontal_pole:Nnn
\coffin_gset_horizontal_pole:cnn

```

New: 2012-07-20
Updated: 2019-01-21

`\coffin_set_horizontal_pole:Nnn <coffin> {\<pole>} {\<offset>}`

Sets the $\langle pole \rangle$ to run horizontally through the $\langle coffin \rangle$. The $\langle pole \rangle$ is placed at the $\langle offset \rangle$ from the bottom edge of the bounding box of the $\langle coffin \rangle$. The $\langle offset \rangle$ should be given as a dimension expression.

```

\coffin_set_vertical_pole:Nnn
\coffin_set_vertical_pole:cnn
\coffin_gset_vertical_pole:Nnn
\coffin_gset_vertical_pole:cnn

```

New: 2012-07-20
Updated: 2019-01-21

`\coffin_set_vertical_pole:Nnn <coffin> {\<pole>} {\<offset>}`

Sets the $\langle pole \rangle$ to run vertically through the $\langle coffin \rangle$. The $\langle pole \rangle$ is placed at the $\langle offset \rangle$ from the left-hand edge of the bounding box of the $\langle coffin \rangle$. The $\langle offset \rangle$ should be given as a dimension expression.

3 Coffin affine transformations

<code>\coffin_resize:Nnn</code>	<code>\coffin_resize:Nnn <coffin> {<width>} {<total-height>}</code>
<code>\coffin_resize:cnn</code>	
<code>\coffin_gresize:Nnn</code>	Resized the $\langle coffin \rangle$ to $\langle width \rangle$ and $\langle total-height \rangle$, both of which should be given as dimension expressions.
<code>\coffin_gresize:cnn</code>	
Updated: 2019-01-23	
<code>\coffin_rotate:Nn</code>	<code>\coffin_rotate:Nn <coffin> {<angle>}</code>
<code>\coffin_rotate:cn</code>	
<code>\coffin_grotate:Nn</code>	Rotates the $\langle coffin \rangle$ by the given $\langle angle \rangle$ (given in degrees counter-clockwise). This process rotates both the coffin content and poles. Multiple rotations do not result in the bounding box of the coffin growing unnecessarily.
<code>\coffin_grotate:cn</code>	
<code>\coffin_scale:Nnn</code>	<code>\coffin_scale:Nnn <coffin> {<x-scale>} {<y-scale>}</code>
<code>\coffin_scale:cnn</code>	
<code>\coffin_gscale:Nnn</code>	Scales the $\langle coffin \rangle$ by a factors $\langle x-scale \rangle$ and $\langle y-scale \rangle$ in the horizontal and vertical directions, respectively. The two scale factors should be given as real numbers.
<code>\coffin_gscale:cnn</code>	
Updated: 2019-01-23	

4 Joining and using coffins

<code>\coffin_attach:NnnNnnnn</code>	<code>\coffin_attach:NnnNnnnn</code>
<code>\coffin_attach:(cnnNnnnn Nnnncnnnn cnnncnnnn)</code>	$\langle coffin_1 \rangle$ { $\langle coffin_1-pole_1 \rangle$ } { $\langle coffin_1-pole_2 \rangle$ }
<code>\coffin_gattach:NnnNnnnn</code>	$\langle coffin_2 \rangle$ { $\langle coffin_2-pole_1 \rangle$ } { $\langle coffin_2-pole_2 \rangle$ }
<code>\coffin_gattach:(cnnNnnnn Nnnncnnnn cnnncnnnn)</code>	{ $\langle x-offset \rangle$ } { $\langle y-offset \rangle$ }
Updated: 2019-01-22	
<p>This function attaches $\langle coffin_2 \rangle$ to $\langle coffin_1 \rangle$ such that the bounding box of $\langle coffin_1 \rangle$ is not altered, <i>i.e.</i> $\langle coffin_2 \rangle$ can protrude outside of the bounding box of the coffin. The alignment is carried out by first calculating $\langle handle_1 \rangle$, the point of intersection of $\langle coffin_1-pole_1 \rangle$ and $\langle coffin_1-pole_2 \rangle$, and $\langle handle_2 \rangle$, the point of intersection of $\langle coffin_2-pole_1 \rangle$ and $\langle coffin_2-pole_2 \rangle$. $\langle coffin_2 \rangle$ is then attached to $\langle coffin_1 \rangle$ such that the relationship between $\langle handle_1 \rangle$ and $\langle handle_2 \rangle$ is described by the $\langle x-offset \rangle$ and $\langle y-offset \rangle$. The two offsets should be given as dimension expressions.</p>	

<code>\coffin_join:NnnNnnnn</code>	<code>\coffin_join:NnnNnnnn</code>
<code>\coffin_join:(cnnNnnnn Nnnncnnnn cnnncnnnn)</code>	$\langle coffin_1 \rangle$ { $\langle coffin_1-pole_1 \rangle$ } { $\langle coffin_1-pole_2 \rangle$ }
<code>\coffin_gjoin:NnnNnnnn</code>	$\langle coffin_2 \rangle$ { $\langle coffin_2-pole_1 \rangle$ } { $\langle coffin_2-pole_2 \rangle$ }
<code>\coffin_gjoin:(cnnNnnnn Nnnncnnnn cnnncnnnn)</code>	{ $\langle x-offset \rangle$ } { $\langle y-offset \rangle$ }
Updated: 2019-01-22	

This function joins $\langle coffin_2 \rangle$ to $\langle coffin_1 \rangle$ such that the bounding box of $\langle coffin_1 \rangle$ may expand. The new bounding box covers the area containing the bounding boxes of the two original coffins. The alignment is carried out by first calculating $\langle handle_1 \rangle$, the point of intersection of $\langle coffin_1-pole_1 \rangle$ and $\langle coffin_1-pole_2 \rangle$, and $\langle handle_2 \rangle$, the point of intersection of $\langle coffin_2-pole_1 \rangle$ and $\langle coffin_2-pole_2 \rangle$. $\langle coffin_2 \rangle$ is then attached to $\langle coffin_1 \rangle$ such that the relationship between $\langle handle_1 \rangle$ and $\langle handle_2 \rangle$ is described by the $\langle x-offset \rangle$ and $\langle y-offset \rangle$. The two offsets should be given as dimension expressions.

```
\coffin_typeset:Nnnnn
\coffin_typeset:cnnnn
```

Updated: 2012-07-20

```
\coffin_typeset:Nnnnn <coffin> {\pole_1} {\pole_2}
{\<x-offset>} {\<y-offset>}
```

Typesetting is carried out by first calculating $\langle handle \rangle$, the point of intersection of $\langle pole_1 \rangle$ and $\langle pole_2 \rangle$. The coffin is then typeset in horizontal mode such that the relationship between the current reference point in the document and the $\langle handle \rangle$ is described by the $\langle x-offset \rangle$ and $\langle y-offset \rangle$. The two offsets should be given as dimension expressions. Typesetting a coffin is therefore analogous to carrying out an alignment where the “parent” coffin is the current insertion point.

5 Measuring coffins

```
\coffin_dp:N
\coffin_dp:c
```

```
\coffin_dp:N <coffin>
```

Calculates the depth (below the baseline) of the $\langle coffin \rangle$ in a form suitable for use in a $\langle dimension expression \rangle$.

```
\coffin_ht:N
\coffin_ht:c
```

```
\coffin_ht:N <coffin>
```

Calculates the height (above the baseline) of the $\langle coffin \rangle$ in a form suitable for use in a $\langle dimension expression \rangle$.

```
\coffin_wd:N
\coffin_wd:c
```

```
\coffin_wd:N <coffin>
```

Calculates the width of the $\langle coffin \rangle$ in a form suitable for use in a $\langle dimension expression \rangle$.

6 Coffin diagnostics

```
\coffin_display_handles:Nn
\coffin_display_handles:cn
```

Updated: 2011-09-02

```
\coffin_display_handles:Nn <coffin> {\color}
```

This function first calculates the intersections between all of the $\langle poles \rangle$ of the $\langle coffin \rangle$ to give a set of $\langle handles \rangle$. It then prints the $\langle coffin \rangle$ at the current location in the source, with the position of the $\langle handles \rangle$ marked on the coffin. The $\langle handles \rangle$ are labelled as part of this process: the locations of the $\langle handles \rangle$ and the labels are both printed in the $\langle color \rangle$ specified.

```
\coffin_mark_handle:Nnnn
\coffin_mark_handle:cnnn
```

Updated: 2011-09-02

```
\coffin_mark_handle:Nnnn <coffin> {\pole_1} {\pole_2} {\color}
```

This function first calculates the $\langle handle \rangle$ for the $\langle coffin \rangle$ as defined by the intersection of $\langle pole_1 \rangle$ and $\langle pole_2 \rangle$. It then marks the position of the $\langle handle \rangle$ on the $\langle coffin \rangle$. The $\langle handle \rangle$ are labelled as part of this process: the location of the $\langle handle \rangle$ and the label are both printed in the $\langle color \rangle$ specified.

```
\coffin_show_structure:N
\coffin_show_structure:c
```

Updated: 2015-08-01

```
\coffin_show_structure:N <coffin>
```

This function shows the structural information about the $\langle coffin \rangle$ in the terminal. The width, height and depth of the typeset material are given, along with the location of all of the poles of the coffin.

Notice that the poles of a coffin are defined by four values: the x and y co-ordinates of a point that the pole passes through and the x - and y -components of a vector denoting the direction of the pole. It is the ratio between the later, rather than the absolute values, which determines the direction of the pole.

`\coffin_log_structure:N`
`\coffin_log_structure:c`

New: 2014-08-22
Updated: 2015-08-01

`\coffin_log_structure:N` $\langle coffin \rangle$

This function writes the structural information about the $\langle coffin \rangle$ in the log file. See also

`\coffin_show_structure:N` which displays the result in the terminal.

7 Constants and variables

`\c_empty_coffin`

A permanently empty coffin.

`\l_tmpa_coffin`
`\l_tmpb_coffin`

New: 2012-06-19

Scratch coffins for local assignment. These are never used by the kernel code, and so are safe for use with any $\text{\LaTeX}3$ -defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.

`\g_tmpa_coffin`
`\g_tmpb_coffin`

New: 2019-01-24

Scratch coffins for global assignment. These are never used by the kernel code, and so are safe for use with any $\text{\LaTeX}3$ -defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.

Part XXXI

The l3color-base package

Color support

This module provides support for color in L^AT_EX3. At present, the material here is mainly intended to support a small number of low-level requirements in other l3kernel modules.

1 Color in boxes

Controlling the color of text in boxes requires a small number of control functions, so that the boxed material uses the color at the point where it is set, rather than where it is used.

<code>\color_group_begin:</code>
<code>\color_group_end:</code>

New: 2011-09-03

<code>\color_group_begin:</code>
<code>...</code>
<code>\color_group_end:</code>

Creates a color group: one used to “trap” color settings. This grouping is built in to for example `\hbox_set:Nn`.

<code>\color_ensure_current:</code>

New: 2011-09-03

<code>\color_ensure_current:</code>

Ensures that material inside a box uses the foreground color at the point where the box is set, rather than that in force when the box is used. This function should usually be used within a `\color_group_begin: ... \color_group_end:` group.

Part XXXII

The l3luatex package: LuaTeX-specific functions

The LuaTeX engine provides access to the Lua programming language, and with it access to the “internals” of TeX. In order to use this within the framework provided here, a family of functions is available. When used with pdfTeX, pTeX, upTeX or XeTeX these raise an error: use `\sys_if_engine luatex:T` to avoid this. Details on using Lua with the LuaTeX engine are given in the LuaTeX manual.

1 Breaking out to Lua

<code>\lua_now:n</code>	★	<code>\lua_now:n</code>	{ <i><token list></i> }
-------------------------	---	-------------------------	-------------------------------

<code>\lua_now:e</code>	★
-------------------------	---

New: 2018-06-18

The *<token list>* is first tokenized by TeX, which includes converting line ends to spaces in the usual TeX manner and which respects currently-applicable TeX category codes. The resulting *<Lua input>* is passed to the Lua interpreter for processing. Each `\lua_now:n` block is treated by Lua as a separate chunk. The Lua interpreter executes the *<Lua input>* immediately, and in an expandable manner.

TeXhackers note: `\lua_now:e` is a macro wrapper around `\directlua:` when LuaTeX is in use two expansions are required to yield the result of the Lua code.

<code>\lua_shipout_e:n</code>	★
-------------------------------	---

<code>\lua_shipout:n</code>	★
-----------------------------	---

New: 2018-06-18

<code>\lua_shipout:n</code>	{ <i><token list></i> }
-----------------------------	-------------------------------

The *<token list>* is first tokenized by TeX, which includes converting line ends to spaces in the usual TeX manner and which respects currently-applicable TeX category codes. The resulting *<Lua input>* is passed to the Lua interpreter when the current page is finalised (*i.e.* at shipout). Each `\lua_shipout:n` block is treated by Lua as a separate chunk. The Lua interpreter will execute the *<Lua input>* during the page-building routine: no TeX expansion of the *<Lua input>* will occur at this stage.

In the case of the `\lua_shipout_e:n` version the input is fully expanded by TeX in an e-type manner during the shipout operation.

TeXhackers note: At a TeX level, the *<Lua input>* is stored as a “whatsit”.

<code>\lua_escape:n</code>	★
----------------------------	---

<code>\lua_escape:e</code>	★
----------------------------	---

New: 2015-06-29

<code>\lua_escape:n</code>	{ <i><token list></i> }
----------------------------	-------------------------------

Converts the *<token list>* such that it can safely be passed to Lua: embedded backslashes, double and single quotes, and newlines and carriage returns are escaped. This is done by prepending an extra token consisting of a backslash with category code 12, and for the line endings, converting them to `\n` and `\r`, respectively.

TeXhackers note: `\lua_escape:e` is a macro wrapper around `\luaescapestring:` when LuaTeX is in use two expansions are required to yield the result of the Lua code.

2 Lua interfaces

As well as interfaces for T_EX, there are a small number of Lua functions provided here.

<hr/> ltx.utils <hr/>	Most public interfaces provided by the module are stored within the <code>ltx.utils</code> table.
<hr/> l3kernel <hr/>	For compatibility reasons, there are also some deprecated interfaces provided in the <code>l3kernel</code> table. These do not return their result as Lua values but instead print them to T _E X.
<hr/> l3kernel.charcat <hr/>	<code>l3kernel.charcat(<charcode>, <catcode>)</code> Constructs a character of <charcode> and <catcode> and returns the result to T _E X.
<hr/> l3kernel.elapsedtime <hr/>	<code>l3kernel.elapsedtime()</code> Returns the CPU time in <scaled seconds> since the start of the T _E X run or since <code>l3kernel.resettimer</code> was issued. This only measures the time used by the CPU, not the real time, e.g., waiting for user input.
<hr/> ltx.utils.filedump l3kernel.filedump <hr/>	<code><dump> = ltx.utils.filedump(<file>, <offset>, <length>)</code> <code>l3kernel.filedump(<file>, <offset>, <length>)</code> Returns the uppercase hexadecimal representation of the content of the <file> read as bytes. If the <length> is given, only this part of the file is returned; similarly, one may specify the <offset> from the start of the file. If the <length> is not given, the entire file is read starting at the <offset>.
<hr/> ltx.utils.filemd5sum l3kernel.filemdfivesum <hr/>	<code><hash> = ltx.utils.filemd5sum(<file>)</code> <code>l3kernel.filemdfivesum(<file>)</code> Returns the MD5 sum of the file contents read as bytes; note that the result will depend on the nature of the line endings used in the file, in contrast to normal T _E X behaviour. If the <file> is not found, nothing is returned with <i>no error raised</i> .
<hr/> ltx.utils.filemoddate l3kernel.filemoddate <hr/>	<code><date> = ltx.utils.filemoddate(<file>)</code> <code>l3kernel.filemoddate(<file>)</code> Returns the date/time of last modification of the <file> in the format <div style="text-align: center;">D: <year><month><day><hour><minute><second><offset></div> where the latter may be Z (UTC) or <plus-minus><hours>'<minutes>'. If the <file> is not found, nothing is returned with <i>no error raised</i> .
<hr/> ltx.utils.filesize l3kernel.filesize <hr/>	<code>size = ltx.utils.filesize(<file>)</code> <code>l3kernel.filesize(<file>)</code> Returns the size of the <file> in bytes. If the <file> is not found, nothing is returned with <i>no error raised</i> .

<u>l3kernel.resettimer</u>	<code>l3kernel.resettimer()</code> Resets the timer used by <code>l3kernel.elapsedtime</code> .
<u>l3kernel.shellescape</u>	<code>l3kernel.shellescape(<i><cmd></i>)</code> Executes the <i><cmd></i> and prints to the log as for pdfTeX.
<u>l3kernel.strcmp</u>	<code>l3kernel.strcmp(<i><str one></i>, <i><str two></i>)</code> Compares the two strings and returns 0 to TeX if the two are identical.

Part XXXIII

The l3unicode package: Unicode support functions

This module provides Unicode-specific functions along with loading data from a range of Unicode Consortium files. At present, it provides no public functions.

Part XXXIV

The l3text package: text processing

1 l3text documentation

This module deals with manipulation of (formatted) text; such material is comprised of a restricted set of token list content. The functions provided here concern conversion of textual content for example in case changing, generation of bookmarks and extraction to tags. All of the major functions operate by expansion. Begin-group and end-group tokens in the $\langle text \rangle$ are normalized and become $\{$ and $\}$, respectively.

1.1 Expanding text

<code>\text_expand:n</code> *	<code>\text_expand:n</code> $\{\langle text \rangle\}$
-------------------------------	--

New: 2020-01-02

Takes user input $\langle text \rangle$ and expands the content. Protected commands (typically formatting) are left in place, and no processing takes place of math mode material (as delimited by pairs given in `\l_text_math_delims_tl` or as the argument to commands listed in `\l_text_math_arg_tl`). Commands which are neither engine- nor L^AT_EX protected are expanded exhaustively. Any commands listed in `\l_text_expand_exclude_tl`, `\l_text_accents_tl` and `\l_text_letterlike_tl` are excluded from expansion.

<code>\text_declare_expand_equivalent:Nn</code>	<code>\text_declare_expand_equivalent:Nn</code> $\langle cmd \rangle$ $\{\langle replacement \rangle\}$
<code>\text_declare_expand_equivalent:cn</code>	

New: 2020-01-22

Declares that the $\langle replacement \rangle$ tokens should be used whenever the $\langle cmd \rangle$ (a single token) is encountered. The $\langle replacement \rangle$ tokens should be expandable.

1.2 Case changing

<code>\text_lowercase:n</code>	★
<code>\text_uppercase:n</code>	★
<code>\text_titlecase:n</code>	★
<code>\text_titlecase_first:n</code>	★
<code>\text_lowercase:nn</code>	★
<code>\text_uppercase:nn</code>	★
<code>\text_titlecase:nn</code>	★
<code>\text_titlecase_first:nn</code>	★

New: 2019-11-20
Updated: 2020-02-24

`\text_uppercase:n` $\{\langle tokens \rangle\}$
`\text_uppercase:nn` $\{\langle language \rangle\} \{\langle tokens \rangle\}$
Takes user input $\langle text \rangle$ first applies `\text_expand`, then transforms the case of character tokens as specified by the function name. The category code of letters are not changed by this process (at least where they can be represented by the engine as a single token: 8-bit engines may require active characters).

Upper- and lowercase have the obvious meanings. Titlecasing may be regarded informally as converting the first character of the $\langle tokens \rangle$ to uppercase and the rest to lowercase. However, the process is more complex than this as there are some situations where a single lowercase character maps to a special form, for example *ij* in Dutch which becomes *IJ*. The `titlecase_first` variant does not attempt any case changing at all after the first letter has been processed.

Importantly, notice that these functions are intended for working with user *text for typesetting*. For case changing programmatic data see the `l3str` module and discussion there of `\str_lowercase:n`, `\str_uppercase:n` and `\str_foldcase:n`.

Case changing does not take place within math mode material so for example

`\text_uppercase:n { Some~text~$y = mx + c$~with~{Braces} }`

becomes

SOME TEXT \$y = mx + c\$ WITH {BRACES}

The arguments of commands listed in `\l_text_case_exclude_arg_tl` are excluded from case changing; the latter are entirely non-textual content (such as labels).

As is generally true for `expl3`, these functions are designed to work with Unicode input only. As such, UTF-8 input is assumed for *all* engines. When used with Xe_{La}TeX or Lua_{TeX} a full range of Unicode transformations are enabled. Specifically, the standard mappings here follow those defined by the [Unicode Consortium](#) in `UnicodeData.txt` and `SpecialCasing.txt`. In the case of 8-bit engines, mappings are provided for characters which can be represented in output typeset using the T1, T2 and LGR font encodings. Thus for example *ä* can be case-changed using pdf_{TeX}. For p_{TeX} only the ASCII range is covered as the engine treats input outside of this range as east Asian.

Language-sensitive conversions are enabled using the $\langle language \rangle$ argument, and follow Unicode Consortium guidelines. Currently, the languages recognised for special handling are as follows.

- Azeri and Turkish (`az` and `tr`). The case pairs I/i-dotless and I-dot/i are activated for these languages. The combining dot mark is removed when lowercasing I-dot and introduced when upper casing i-dotless.
- German (`de-alt`). An alternative mapping for German in which the lowercase *Eszett* maps to a *großes Eszett*. Since there is a T1 slot for the *großes Eszett* in T1, this tailoring *is* available with pdf_{TeX} as well as in the Unicode _{TeX} engines.
- Greek (`el`). Removes accents from Greek letters when uppercasing; titlecasing leaves accents in place. (At present this is implemented only for Unicode engines.)
- Lithuanian (`lt`). The lowercase letters i and j should retain a dot above when the accents grave, acute or tilde are present. This is implemented for lowercasing of the relevant uppercase letters both when input as single Unicode codepoints and when using combining accents. The combining dot is removed when uppercasing in these cases. Note that *only* the accents used in Lithuanian are covered: the behaviour of other accents are not modified.
- Dutch (`nl`). Capitalisation of *ij* at the beginning of titlecased input produces *IJ* rather than *Ij*. The output retains two separate letters, thus this transformation *is* available using pdf_{TeX}.

For titlecasing, note that there are two functions available. The function `\text_`
`titlecase:n` applies (broadly) uppercasing to the first letter of the input, then lower-

1.3 Removing formatting from text

<code>\text_purify:n</code> *	<code>\text_purify:n {<text>}</code>
-------------------------------	--

New: 2020-03-05
Updated: 2020-05-14

Takes user input $\langle text \rangle$ and expands as described for `\text_expand:n`, then removes all functions from the resulting text. Math mode material (as delimited by pairs given in `\l_text_math_delims_tl` or as the argument to commands listed in `\l_text_math_arg_tl`) is left contained in a pair of \$ delimiters. Non-expandable functions present in the $\langle text \rangle$ must either have a defined equivalent (see `\text_declare_purify_equivalent:Nn`) or will be removed from the result. Implicit tokens are converted to their explicit equivalent.

<code>\text_declare_purify_equivalent:Nn</code>	<code>\text_declare_purify_equivalent:Nn <cmd> {<replacement>}</code>
<code>\text_declare_purify_equivalent:Nx</code>	

New: 2020-03-05

Declares that the $\langle replacement \rangle$ tokens should be used whenever the $\langle cmd \rangle$ (a single token) is encountered. The $\langle replacement \rangle$ tokens should be expandable.

1.4 Control variables

<code>\l_text_accents_tl</code>

Lists commands which represent accents, and which are left unchanged by expansion. (Defined only for the L^AT_EX 2_ε package.)

<code>\l_text_letterlike_tl</code>

Lists commands which represent letters; these are left unchanged by expansion. (Defined only for the L^AT_EX 2_ε package.)

<code>\l_text_math_arg_tl</code>

Lists commands present in the $\langle text \rangle$ where the argument of the command should be treated as math mode material. The treatment here is similar to `\l_text_math_delims_tl` but for a command rather than paired delimiters.

<code>\l_text_math_delims_tl</code>

Lists pairs of tokens which delimit (in-line) math mode content; such content *may* be excluded from processing.

<code>\l_text_case_exclude_arg_tl</code>
--

Lists commands which are excluded from case changing.

<code>\l_text_expand_exclude_tl</code>
--

Lists commands which are excluded from expansion.

<code>\l_text_titlecase_check_letter_bool</code>
--

Controls how the start of titlecasing is handled: when **true**, the first *letter* in text is considered. The standard setting is **true**.

Part XXXV

The l3legacy package

Interfaces to legacy concepts

There are a small number of T_EX or L^AT_EX 2_ε concepts which are not used in `expl3` code but which need to be manipulated when working as a L^AT_EX 2_ε package. To allow these to be integrated cleanly into `expl3` code, a set of legacy interfaces are provided here.

<code>\legacy_if_p:n</code> *	<code>\legacy_if:nTF</code> { <i><name></i> } { <i><true code></i> } { <i><false code></i> }
<code>\legacy_if:nTF</code> *	Tests if the L ^A T _E X 2 _ε /plain T _E X conditional (generated by <code>\newif</code>) if <code>true</code> or <code>false</code> and branches accordingly. The <i><name></i> of the conditional should <i>omit</i> the leading <code>if</code> .

Part XXXVI

The l3candidates package

Experimental additions to l3kernel

1 Important notice

This module provides a space in which functions can be added to l3kernel (expl3) while still being experimental.

As such, the functions here may not remain in their current form, or indeed at all, in l3kernel in the future.

In contrast to the material in l3experimental, the functions here are all *small* additions to the kernel. We encourage programmers to test them out and report back on the LaTeX-L mailing list.

Thus, if you intend to use any of these functions from the candidate module in a public package offered to others for productive use (e.g., being placed on CTAN) please consider the following points carefully:

- Be prepared that your public packages might require updating when such functions are being finalized.
- Consider informing us that you use a particular function in your public package, e.g., by discussing this on the LaTeX-L mailing list. This way it becomes easier to coordinate any updates necessary without issues for the users of your package.
- Discussing and understanding use cases for a particular addition or concept also helps to ensure that we provide the right interfaces in the final version so please give us feedback if you consider a certain candidate function useful (or not).

We only add functions in this space if we consider them being serious candidates for a final inclusion into the kernel. However, real use sometimes leads to better ideas, so functions from this module are **not necessarily stable** and we may have to adjust them!

2 Additions to l3box

2.1 Viewing part of a box

```
\box_clip:N  
\box_clip:c  
\box_gclip:N  
\box_gclip:c
```

Updated: 2019-01-23

```
\box_clip:N <box>
```

Clips the $\langle box \rangle$ in the output so that only material inside the bounding box is displayed in the output. The updated $\langle box \rangle$ is an hbox, irrespective of the nature of the $\langle box \rangle$ before the clipping is applied.

These functions require the L^AT_EX3 native drivers: they do not work with the L^AT_EX 2_ε graphics drivers!

T_EXhackers note: Clipping is implemented by the driver, and as such the full content of the box is placed in the output file. Thus clipping does not remove any information from the raw output, and hidden material can therefore be viewed by direct examination of the file.

```
\box_set_trim:Nnnnn
\box_set_trim:cnnnn
\box_gset_trim:Nnnnn
\box_gset_trim:cnnnn
```

New: 2019-01-23

```
\box_set_trim:Nnnnn <box> {\left} {\bottom} {\right} {\top}
```

Adjusts the bounding box of the $\langle box \rangle$ $\langle left \rangle$ is removed from the left-hand edge of the bounding box, $\langle right \rangle$ from the right-hand edge and so fourth. All adjustments are *$\langle dimension expressions \rangle$* . Material outside of the bounding box is still displayed in the output unless `\box_clip:N` is subsequently applied. The updated $\langle box \rangle$ is an hbox, irrespective of the nature of the $\langle box \rangle$ before the trim operation is applied. The behavior of the operation where the trims requested is greater than the size of the box is undefined.

```
\box_set_viewport:Nnnnn
\box_set_viewport:cnnnn
\box_gset_viewport:Nnnnn
\box_gset_viewport:cnnnn
```

New: 2019-01-23

```
\box_set_viewport:Nnnnn <box> {\llx} {\lly} {\urx} {\ury}
```

Adjusts the bounding box of the $\langle box \rangle$ such that it has lower-left co-ordinates ($\langle llx \rangle$, $\langle lly \rangle$) and upper-right co-ordinates ($\langle urx \rangle$, $\langle ury \rangle$). All four co-ordinate positions are *$\langle dimension expressions \rangle$* . Material outside of the bounding box is still displayed in the output unless `\box_clip:N` is subsequently applied. The updated $\langle box \rangle$ is an hbox, irrespective of the nature of the $\langle box \rangle$ before the viewport operation is applied.

3 Additions to l3expan

```
\exp_args_generate:n
```

New: 2018-04-04
Updated: 2019-02-08

```
\exp_args_generate:n {\variant argument specifiers}
```

Defines `\exp_args:N<variant>` functions for each $\langle variant \rangle$ given in the comma list $\{\langle variant argument specifiers \rangle\}$. Each $\langle variant \rangle$ should consist of the letters N, c, n, V, v, o, f, e, x, p and the resulting function is protected if the letter x appears in the $\langle variant \rangle$. This is only useful for cases where `\cs_generate_variant:Nn` is not applicable.

4 Additions to l3fp

```
\fp_if_nan_p:n ★
\fp_if_nan:nTF ★
```

New: 2019-08-25

```
\fp_if_nan:n {\fpexpr}
```

Evaluates the $\langle fpexpr \rangle$ and tests whether the result is exactly NaN. The test returns **false** for any other result, even a tuple containing NaN.

5 Additions to l3file

```
\iow_allow_break:
```

New: 2018-12-29

```
\iow_allow_break:
```

In the first argument of `\iow_wrap:nnnn` (for instance in messages), inserts a break-point that allows a line break. In other words this is a zero-width breaking space.

`\ior_get_term:nN`
`\ior_str_get_term:nN`

New: 2019-03-23

`\ior_get_term:nN` $\langle prompt \rangle$ $\langle token\ list\ variable \rangle$

Function that reads one or more lines (until an equal number of left and right braces are found) from the terminal and stores the result locally in the $\langle token\ list \rangle$ variable. Tokenization occurs as described for `\ior_get:NN` or `\ior_str_get:NN`, respectively. When the $\langle prompt \rangle$ is empty, \TeX will wait for input without any other indication: typically the programmer will have provided a suitable text using e.g. `\iow_term:n`. Where the $\langle prompt \rangle$ is given, it will appear in the terminal followed by an =, e.g.

prompt=

`\ior_shell_open:Nn`

New: 2019-05-08

`\ior_shell_open:Nn` $\langle stream \rangle$ $\{ \langle shell\ command \rangle \}$

Opens the *pseudo*-file created by the output of the $\langle shell\ command \rangle$ for reading using $\langle stream \rangle$ as the control sequence for access. If the $\langle stream \rangle$ was already open it is closed before the new operation begins. The $\langle stream \rangle$ is available for access immediately and will remain allocated to $\langle shell\ command \rangle$ until a `\ior_close:N` instruction is given or the \TeX run ends. If piped system calls are disabled an error is raised.

For details of handling of the $\langle shell\ command \rangle$, see `\sys_get_shell:nn(TF)`.

6 Additions to `\l3flag`

`\flag_raise_if_clear:n` ☆

New: 2018-04-02

`\flag_raise_if_clear:n` $\{ \langle flag\ name \rangle \}$

Ensures the $\langle flag \rangle$ is raised by making its height at least 1, locally.

7 Additions to `\l3intarray`

`\intarray_gset_rand:Nnn`
`\intarray_gset_rand:cnn`
`\intarray_gset_rand:Nn`
`\intarray_gset_rand:cn`

New: 2018-05-05

`\intarray_gset_rand:Nnn` $\langle intarray\ var \rangle$ $\{ \langle minimum \rangle \}$ $\{ \langle maximum \rangle \}$
`\intarray_gset_rand:Nn` $\langle intarray\ var \rangle$ $\{ \langle maximum \rangle \}$

Evaluates the integer expressions $\langle minimum \rangle$ and $\langle maximum \rangle$ then sets each entry (independently) of the $\langle integer\ array\ variable \rangle$ to a pseudo-random number between the two (with bounds included). If the absolute value of either bound is bigger than $2^{30} - 1$, an error occurs. Entries are generated in the same way as repeated calls to `\int_rand:nn` or `\int_rand:n` respectively, in particular for the second function the $\langle minimum \rangle$ is 1. Assignments are always global. This is not available in older versions of \LaTeX .

7.1 Working with contents of integer arrays

`\intarray_to_clist:N` ☆

New: 2018-05-04

`\intarray_to_clist:N` $\langle intarray\ var \rangle$

Converts the $\langle intarray \rangle$ to integer denotations separated by commas. All tokens have category code other. If the $\langle intarray \rangle$ has no entry the result is empty; otherwise the result has one fewer comma than the number of items.

8 Additions to l3msg

```
\msg_show_eval:Nn
\msg_log_eval:Nn
```

New: 2017-12-04

```
\msg_show_eval:Nn <function> {<expression>}
```

Shows or logs the $\langle expression \rangle$ (turned into a string), an equal sign, and the result of applying the $\langle function \rangle$ to the $\{ \langle expression \rangle \}$ (with f-expansion). For instance, if the $\langle function \rangle$ is `\int_eval:n` and the $\langle expression \rangle$ is `1+2` then this logs `> 1+2=3.`

```
\msg_show:nnnnnn
\msg_show:nnxxxx
\msg_show:nnnnn
\msg_show:nnxxx
\msg_show:nnnn
\msg_show:nnxx
\msg_show:nnn
\msg_show:nnx
\msg_show:nn
```

New: 2017-12-04

```
\msg_show:nnnnnn {<module>} {<message>} {<arg one>} {<arg two>} {<arg three>} {<arg four>}
```

Issues $\langle module \rangle$ information $\langle message \rangle$, passing $\langle arg one \rangle$ to $\langle arg four \rangle$ to the text-creating functions. The information text is shown on the terminal and the TeX run is interrupted in a manner similar to `\tl_show:n`. This is used in conjunction with `\msg_show_item:n` and similar functions to print complex variable contents completely. If the formatted text does not contain `>~` at the start of a line, an additional line `>~.` will be put at the end. In addition, a final period is added if not present.

```
\msg_show_item:n          * \seq_map_function:NN <seq> \msg_show_item:n
\msg_show_item_unbraced:n * \prop_map_function:NN <prop> \msg_show_item:nn
\msg_show_item:nn         *
\msg_show_item_unbraced:nn *
```

New: 2017-12-04

Used in the text of messages for `\msg_show:nnxxxx` to show or log a list of items or key-value pairs. The one-argument functions are used for sequences, clist or token lists and the others for property lists. These functions turn their arguments to strings.

9 Additions to l3prg

```
\bool_set_inverse:N
\bool_set_inverse:c
\bool_gset_inverse:N
\bool_gset_inverse:c
```

New: 2018-05-10

```
\bool_set_inverse:N <boolean>
```

Toggles the $\langle boolean \rangle$ from true to false and conversely: sets it to the inverse of its current value.

<hr/>	
<code>\bool_case_true:n</code>	★
<code>\bool_case_true:nTF</code>	★
<code>\bool_case_false:n</code>	★
<code>\bool_case_false:nTF</code>	★
<hr/>	
New: 2019-02-10	
<hr/>	
<pre> \bool_case_true:nTF { {<boolexpr case₁>} {<code case₁>} {<boolexpr case₂>} {<code case₂>} ... {<boolexpr case_n>} {<code case_n>} } {<true code>} {<false code>} </pre>	

Evaluates in turn each of the *<boolean expression cases>* until the first one that evaluates to true or to false, for `\bool_case_true:n` and `\bool_case_false:n`, respectively. The *<code>* associated to this first case is left in the input stream, followed by the *<true code>*, and other cases are discarded. If none of the cases match then only the *<false code>* is inserted. The functions `\bool_case_true:n` and `\bool_case_false:n`, which do nothing if there is no match, are also available. For example

```

\bool_case_true:nF
{
  { \dim_compare_p:n { \l__mypkg_wd_dim <= 10pt } }
    { Fits }
  { \int_compare_p:n { \l__mypkg_total_int >= 10 } }
    { Many }
  { \l__mypkg_special_bool }
    { Special }
}
{ No idea! }

```

leaves “Fits” or “Many” or “Special” or “No idea!” in the input stream, in a way similar to some other language’s “if ... elseif ... elseif ... else ...”.

10 Additions to l3prop

<hr/>	
<code>\prop_rand_key_value:N</code>	★
<code>\prop_rand_key_value:c</code>	★
<hr/>	
New: 2016-12-06	
<hr/>	
<pre> \prop_rand_key_value:N <prop var> </pre>	

Selects a pseudo-random key–value pair from the *<property list>* and returns *{<key>}* and *{<value>}*. If the *<property list>* is empty the result is empty. This is not available in older versions of X_YTeX.

TeXhackers note: The result is returned within the `\unexpanded` primitive (`\exp_not:n`), which means that the *<value>* does not expand further when appearing in an *x*-type argument expansion.

11 Additions to l3seq

<code>\seq_mapthread_function:NNN</code>	☆	<code>\seq_mapthread_function:NNN <seq₁> <seq₂> <function></code>
<code>\seq_mapthread_function:(NcN cNN ccN)</code>	☆	

Applies *<function>* to every pair of items *<seq₁-item>*–*<seq₂-item>* from the two sequences, returning items from both sequences from left to right. The *<function>* receives two *n*-type arguments for each iteration. The mapping terminates when the end of either sequence is reached (*i.e.* whichever sequence has fewer items determines how many iterations occur).

<code>\seq_set_filter:NNn</code>	<code>\seq_set_filter:NNn <sequence₁> <sequence₂> {<inline boolexpr>}</code>
<code>\seq_gset_filter:NNn</code>	

Evaluates the *<inline boolexpr>* for every *<item>* stored within the *<sequence₂>*. The *<inline boolexpr>* receives the *<item>* as #1. The sequence of all *<items>* for which the *<inline boolexpr>* evaluated to `true` is assigned to *<sequence₁>*.

TeXhackers note: Contrarily to other mapping functions, `\seq_map_break:` cannot be used in this function, and would lead to low-level TeX errors.

<code>\seq_set_from_function:NnN</code>	<code>\seq_set_from_function:NnN <seq var> {<loop code>} <function></code>
<code>\seq_gset_from_function:NnN</code>	

New: 2018-04-06

Sets the *<seq var>* equal to a sequence whose items are obtained by *x*-expanding *<loop code>* *<function>*. This expansion must result in successive calls to the *<function>* with no nonexpandable tokens in between. More precisely the *<function>* is replaced by a wrapper function that inserts the appropriate separators between items in the sequence. The *<loop code>* must be expandable; it can be for example `\tl_map_function:NN <tl var>` or `\clist_map_function:nN {<clist>}` or `\int_step_function:nnnN {<initial value>} {<step>} {<final value>}`.

<code>\seq_set_from_inline_x:Nnn</code>	<code>\seq_set_from_inline_x:Nnn <seq var> {<loop code>} {<inline code>}</code>
<code>\seq_gset_from_inline_x:Nnn</code>	

New: 2018-04-06

Sets the *<seq var>* equal to a sequence whose items are obtained by *x*-expanding *<loop code>* applied to a *<function>* derived from the *<inline code>*. A *<function>* is defined, that takes one argument, *x*-expands the *<inline code>* with that argument as #1, then adds appropriate separators to turn the result into an item of the sequence. The *x*-expansion of *<loop code>* *<function>* must result in successive calls to the *<function>* with no nonexpandable tokens in between. The *<loop code>* must be expandable; it can be for example `\tl_map_function:NN <tl var>` or `\clist_map_function:nN {<clist>}` or `\int_step_function:nnnN {<initial value>} {<step>} {<final value>}`, but not the analogous “inline” mappings.

12 Additions to l3sys

`\c_sys_engine_version_str`

New: 2018-05-02

The version string of the current engine, in the same form as given in the banner issued when running a job. For pdfTeX and LuaTeX this is of the form

$$\langle major \rangle . \langle minor \rangle . \langle revision \rangle$$

For XeTeX, the form is

$$\langle major \rangle . \langle minor \rangle$$

For pTeX and upTeX, only releases since TeX Live 2018 make the data available, and the form is more complex, as it comprises the pTeX version, the upTeX version and the e-pTeX version.

$$p \langle major \rangle . \langle minor \rangle . \langle revision \rangle - u \langle major \rangle . \langle minor \rangle - \langle epTeX \rangle$$

where the u part is only present for upTeX.

`\sys_if_rand_exist_p: *`
`\sys_if_rand_exist:TF *`

New: 2017-05-27

`\sys_if_rand_exist_p:`
`\sys_if_rand_exist:TF {<true code>} {<false code>}`

Tests if the engine has a pseudo-random number generator. Currently this is the case in pdfTeX, LuaTeX, pTeX, upTeX and recent releases of XeTeX.

13 Additions to l3tl

<code>\tl_range_braced:Nnn</code>	★	<code>\tl_range_braced:Nnn <tl var> {<start index>} {<end index>}</code>
<code>\tl_range_braced:cnn</code>	★	<code>\tl_range_braced:nnn {<token list>} {<start index>} {<end index>}</code>
<code>\tl_range_braced:nnn</code>	★	<code>\tl_range_unbraced:Nnn <tl var> {<start index>} {<end index>}</code>
<code>\tl_range_unbraced:Nnn</code>	★	<code>\tl_range_unbraced:nnn {<token list>} {<start index>} {<end index>}</code>
<code>\tl_range_unbraced:cnn</code>	★	Leaves in the input stream the items from the <i><start index></i> to the <i><end index></i> inclusive, using the same indexing as <code>\tl_range:nnn</code> . Spaces are ignored. Regardless of whether items appear with or without braces in the <i><token list></i> , the <code>\tl_range_braced:nnn</code> function wraps each item in braces, while <code>\tl_range_unbraced:nnn</code> does not (overall it removes an outer set of braces). For instance,
<code>\tl_range_unbraced:nnn</code>	★	

New: 2017-07-15

```

\tl_iow_term:x { \tl_range_braced:nnn { abcd~{e}}f } { 2 } { 5 } }
\tl_iow_term:x { \tl_range_braced:nnn { abcd~{e}}f } { -4 } { -1 } }
\tl_iow_term:x { \tl_range_braced:nnn { abcd~{e}}f } { -2 } { -1 } }
\tl_iow_term:x { \tl_range_braced:nnn { abcd~{e}}f } { 0 } { -1 } }

```

prints `{b}{c}{d}{e}}`, `{c}{d}{e}}{f}`, `{e}}{f}`, and an empty line to the terminal, while

```

\tl_iow_term:x { \tl_range_unbraced:nnn { abcd~{e}}f } { 2 } { 5 } }
\tl_iow_term:x { \tl_range_unbraced:nnn { abcd~{e}}f } { -4 } { -1 } }
\tl_iow_term:x { \tl_range_unbraced:nnn { abcd~{e}}f } { -2 } { -1 } }
\tl_iow_term:x { \tl_range_unbraced:nnn { abcd~{e}}f } { 0 } { -1 } }

```

prints `bcde{}`, `cde{f}`, `e{f}`, and an empty line to the terminal. Because braces are removed, the result of `\tl_range_unbraced:nnn` may have a different number of items as for `\tl_range:nnn` or `\tl_range_braced:nnn`. In cases where preserving spaces is important, consider the slower function `\tl_range:nnn`.

T_EXhackers note: The result is returned within the `\unexpanded` primitive (`\exp_not:n`), which means that the *<item>* does not expand further when appearing in an x-type argument expansion.

<code>\tl_build_begin:N</code>	<code>\tl_build_begin:N <tl var></code>
<code>\tl_build_gbegin:N</code>	Clears the <i><tl var></i> and sets it up to support other <code>\tl_build...</code> functions, which allow accumulating large numbers of tokens piece by piece much more efficiently than standard l3tl functions. Until <code>\tl_build_end:N <tl var></code> is called, applying any function from l3tl other than <code>\tl_build...</code> will lead to incorrect results. The <code>begin</code> and <code>gbegin</code> functions must be used for local and global <i><tl var></i> respectively.

New: 2018-04-01

<code>\tl_build_clear:N</code>	<code>\tl_build_clear:N <tl var></code>
<code>\tl_build_gclear:N</code>	Clears the <i><tl var></i> and sets it up to support other <code>\tl_build...</code> functions. The <code>clear</code> and <code>gclear</code> functions must be used for local and global <i><tl var></i> respectively.

New: 2018-04-01

```

\tl_build_put_left:Nn
\tl_build_put_left:Nx
\tl_build_gput_left:Nn
\tl_build_gput_left:Nx
\tl_build_put_right:Nn
\tl_build_put_right:Nx
\tl_build_gput_right:Nn
\tl_build_gput_right:Nx

```

New: 2018-04-01

```

\tl_build_get:NN

```

New: 2018-04-01

```

\tl_build_end:N
\tl_build_gend:N

```

New: 2018-04-01

```

\tl_build_put_left:Nn <tl var> {<tokens>}
\tl_build_put_right:Nn <tl var> {<tokens>}

```

Adds *<tokens>* to the left or right side of the current contents of *<tl var>*. The *<tl var>* must have been set up with `\tl_build_begin:N` or `\tl_build_gbegin:N`. The `put` and `gput` functions must be used for local and global *<tl var>* respectively. The `right` functions are about twice faster than the `left` functions.

```

\tl_build_get:N <tl var1> <tl var2>

```

Stores the contents of the *<tl var₁>* in the *<tl var₂>*. The *<tl var₁>* must have been set up with `\tl_build_begin:N` or `\tl_build_gbegin:N`. The *<tl var₂>* is a “normal” token list variable, assigned locally using `\tl_set:Nn`.

```

\tl_build_end:N <tl var>

```

Gets the contents of *<tl var>* and stores that into the *<tl var>* using `\tl_set:Nn` or `\tl_gset:Nn`. The *<tl var>* must have been set up with `\tl_build_begin:N` or `\tl_build_gbegin:N`. The `end` and `gend` functions must be used for local and global *<tl var>* respectively. These functions completely remove the setup code that enabled *<tl var>* to be used for other `\tl_build_...` functions.

14 Additions to l3token

```

\c_catcode_active_space_tl

```

New: 2017-08-07

```

\char_to_utfviii_bytes:n ★

```

New: 2020-01-09

Token list containing one character with category code 13, (“active”), and character code 32 (space).

```

\char_to_utfviii_bytes:n {<codepoint>}

```

Converts the (Unicode) *<codepoint>* to UTF-8 bytes. The expansion of this function comprises four brace groups, each of which will contain a hexadecimal value: the appropriate byte. As UTF-8 is a variable-length, one or more of the groups may be empty: the bytes read in the logical order, such that a two-byte codepoint will have groups **#1** and **#2** filled and **#3** and **#4** empty.

```

\char_to_nfcd:N ☆

```

New: 2020-01-02

```

\char_to_nfcd:N <char>

```

Converts the *<char>* to the Unicode Normalization Form Canonical Decomposition. The category code of the generated character is the same as the *<char>*. With 8-bit engines, no change is made to the character.

<code>\peek_catcode_collect_inline:Nn</code>	<code>\peek_catcode_collect_inline:Nn <test token> {<inline code>}</code>
<code>\peek_charcode_collect_inline:Nn</code>	<code>\peek_charcode_collect_inline:Nn <test token> {<inline code>}</code>
<code>\peek_meaning_collect_inline:Nn</code>	<code>\peek_meaning_collect_inline:Nn <test token> {<inline code>}</code>

New: 2018-09-23

Collects and removes tokens from the input stream until finding a token that does not match the `<test token>` (as defined by the test `\token_if_eq_catcode:NNTF` or `\token_if_eq_charcode:NNTF` or `\token_if_eq_meaning:NNTF`). The collected tokens are passed to the `<inline code>` as #1. When begin-group or end-group tokens (usually `{` or `}`) are collected they are replaced by implicit `\c_group_begin_token` and `\c_group_end_token`, and when spaces (including `\c_space_token`) are collected they are replaced by explicit spaces.

For example the following code prints “Hello” to the terminal and leave “, world!” in the input stream.

```
\peek_catcode_collect_inline:Nn A { \iow_term:n {#1} } Hello,~world!
```

Another example is that the following code tests if the next token is `*`, ignoring intervening spaces, but putting them back using `#1` if there is no `*`.

```
\peek_meaning_collect_inline:Nn \c_space_token
{ \peek_charcode:NNTF * { star } { no~star #1 } }
```

<code>\peek_remove_spaces:n</code>	<code>\peek_remove_spaces:n {<code>}</code>
------------------------------------	---

New: 2018-10-01

Removes explicit and implicit space tokens (category code 10 and character code 32) from the input stream, then inserts `<code>`.

Part XXXVII

Implementation

1 l3bootstrap implementation

```
1 <*package>
2 <@@=kernel>
```

1.1 LuaTeX-specific code

Depending on the versions available, the L^AT_EX format may not have the raw `\Umath` primitive names available. We fix that globally: it should cause no issues. Older LuaTeX versions do not have a pre-built table of the primitive names here so sort one out ourselves. These end up globally-defined but at that is true with a newer format anyway and as they all start `\U` this should be reasonably safe.

```
3 \begingroup
4   \expandafter\ifx\csname directlua\endcsname\relax
5   \else
6     \directlua{%
7       local i
8       local t = { }
```

```

9      for _,i in pairs(tex.extraprimatives("luatex")) do
10         if string.match(i,"^U") then
11            if not string.match(i,"^Uchar$") then %$
12               table.insert(t,i)
13            end
14         end
15      end
16      tex.enableprimitives("", t)
17   }%
18   \fi
19 \endgroup

```

1.2 The `\pdfstrcmp` primitive in \XeTeX

Only \pdfTeX has a primitive called `\pdfstrcmp`. The \XeTeX version is just `\strcmp`, so there is some shuffling to do. As this is still a real primitive, using the \pdfTeX name is “safe”.

```

20 \begingroup\expandafter\expandafter\expandafter\endgroup
21 \expandafter\ifx\csname pdfstrcmp\endcsname\relax
22 \let\pdfstrcmp\strcmp
23 \fi

```

1.3 Loading support Lua code

When \LuaTeX is used there are various pieces of Lua code which need to be loaded. The code itself is defined in `l3luatex` and is extracted into a separate file. Thus here the task is to load the Lua code both now and (if required) at the start of each job.

```

24 \begingroup\expandafter\expandafter\expandafter\endgroup
25 \expandafter\ifx\csname directlua\endcsname\relax
26 \else
27 \ifnum\luatexversion<110 %
28 \else

```

For \LuaTeX we make sure the basic support is loaded: this is only necessary in plain.

Additionally we just ensure that \TeX has seen the csnames `\prg_return_true:` and `\prg_return_false:` before the Lua code builds these tokens.

```

29 \begingroup\expandafter\expandafter\expandafter\endgroup
30 \expandafter\ifx\csname newcatcodetable\endcsname\relax
31 \input{ltluatex}%
32 \fi
33 \begingroup\edef\ignored{%
34 \expandafter\noexpand\csname prg_return_true:\endcsname
35 \expandafter\noexpand\csname prg_return_false:\endcsname
36 }\endgroup
37 \directlua{require("expl3")}%

```

As the user might be making a custom format, no assumption is made about matching package mode with only loading the Lua code once. Instead, a query to Lua reveals what mode is in operation.

```

38 \ifnum 0%
39 \directlua{
40 if status.ini_version then
41 tex.write("1")

```

```

42     end
43   }>0 %
44   \everyjob\expandafter{%
45     \the\expandafter\everyjob
46     \csname\detokenize{lua_now:n}\endcsname{require("expl3")}%
47   }%
48 \fi
49 \fi
50 \fi

```

1.4 Engine requirements

The code currently requires ε -TeX and functionality equivalent to `\pdfstrcmp`, and also driver and Unicode character support. This is available in a reasonably-wide range of engines.

For LuaTeX, we require at least Lua 5.3 and the `token.set_lua` function. This is available at least since LuaTeX 1.10.

```

51 \begingroup
52 \def\next{\endgroup}%
53 \def\ShortText{Required primitives not found}%
54 \def\LongText%
55   {%
56     LaTeX3 requires the e-TeX primitives and additional functionality as
57     described in the README file.
58     \LineBreak
59     These are available in the engines\LineBreak
60     - pdfTeX v1.40\LineBreak
61     - XeTeX v0.99992\LineBreak
62     - LuaTeX v1.10\LineBreak
63     - e-(u)pTeX mid-2012\LineBreak
64     or later.\LineBreak
65     \LineBreak
66   }%
67 \ifnum0%
68   \expandafter\ifx\csname pdfstrcmp\endcsname\relax
69   \else
70     \expandafter\ifx\csname pdftexversion\endcsname\relax
71       \expandafter\ifx\csname Ucharcat\endcsname\relax
72         \expandafter\ifx\csname kanjiskip\endcsname\relax
73         \else
74           1%
75         \fi
76       \else
77         1%
78       \fi
79     \else
80       \ifnum\pdftexversion<140 \else 1\fi
81     \fi
82   \fi
83   \expandafter\ifx\csname directlua\endcsname\relax
84   \else
85     \ifnum\luatexversion<110 \else 1\fi
86   \fi

```



```

87      =0 %
88      \newlinechar‘^^J %
89      \def\LineBreak{\noexpand\MessageBreak}%
90      \expandafter\ifx\csname PackageError\endcsname\relax
91      \def\LineBreak{^^J}%
92      \def\PackageError#1#2#3%
93      {%
94      \errhelp{#3}%
95      \errmessage{#1 Error: #2}%
96      }%
97      \fi
98      \edef\next
99      {%
100      \noexpand\PackageError{expl3}{\ShortText}
101      {\LongText Loading of expl3 will abort!}%
102      \endgroup
103      \noexpand\endinput
104      }%
105      \fi
106      \next

```

1.5 Extending allocators

The ability to extend \TeX ’s allocation routine to allow for $\varepsilon\text{-}\text{\TeX}$ has been around since 1997 in the `etex` package. Loading this support is delayed until here as we are now sure that the $\varepsilon\text{-}\text{\TeX}$ extensions and `\pdfstrcmp` or equivalent are available. Thus there is no danger of an “uncontrolled” error if the engine requirements are not met.

For $\text{\LaTeX}2_{\varepsilon}$ we need to make sure that the extended pool is being used: `expl3` uses a lot of registers. For formats from 2015 onward there is nothing to do as this is automatic. For older formats, the `etex` package needs to be loaded to do the job. In that case, some inserts are reserved also as these have to be from the standard pool. Note that `\reserveinserts` is `\outer` and so is accessed here by `csname`. In earlier versions, loading `etex` was done directly and so `\reserveinserts` appeared in the code: this then required a `\relax` after `\RequirePackage` to prevent an error with “unsafe” definitions as seen for example with `capoptions`. The optional loading here is done using a group and `\ifx` test as we are not quite in the position to have a single name for `\pdfstrcmp` just yet.

```

107 \begingroup
108 \def\@tempa{LaTeX2e}%
109 \def\next{}%
110 \ifx\fmtname\@tempa
111 \expandafter\ifx\csname extrafloats\endcsname\relax
112 \def\next
113 {%
114 \RequirePackage{etex}%
115 \csname reserveinserts\endcsname{32}%
116 }%
117 \fi
118 \fi
119 \expandafter\endgroup
120 \next

```

1.6 The L^AT_EX3 code environment

The code environment is now set up.

\ExplSyntaxOff Before changing any category codes, in package mode we need to save the situation before loading. Note the set up here means that once applied `\ExplSyntaxOff` becomes a “do nothing” command until `\ExplSyntaxOn` is used.

```

121 \protected\edef\ExplSyntaxOff
122   {%
123     \protected\def\noexpand\ExplSyntaxOff{}%
124     \catcode 9 = \the\catcode 9\relax
125     \catcode 32 = \the\catcode 32\relax
126     \catcode 34 = \the\catcode 34\relax
127     \catcode 38 = \the\catcode 38\relax
128     \catcode 58 = \the\catcode 58\relax
129     \catcode 94 = \the\catcode 94\relax
130     \catcode 95 = \the\catcode 95\relax
131     \catcode 124 = \the\catcode 124\relax
132     \catcode 126 = \the\catcode 126\relax
133     \endlinechar = \the\endlinechar\relax
134     \chardef\csname\detokenize{l__kernel_expl_bool}\endcsname = 0\relax
135   }%

```

(End definition for `\ExplSyntaxOff`. This function is documented on page 7.)

The code environment is now set up.

```

136 \catcode 9 = 9\relax
137 \catcode 32 = 9\relax
138 \catcode 34 = 12\relax
139 \catcode 38 = 4\relax
140 \catcode 58 = 11\relax
141 \catcode 94 = 7\relax
142 \catcode 95 = 11\relax
143 \catcode 124 = 12\relax
144 \catcode 126 = 10\relax
145 \endlinechar = 32\relax

```

`\l__kernel_expl_bool` The status for experimental code syntax: this is on at present.

```

146 \chardef\l__kernel_expl_bool = 1\relax

```

(End definition for `\l__kernel_expl_bool`.)

\ExplSyntaxOn The idea here is that multiple `\ExplSyntaxOn` calls are not going to mess up category codes, and that multiple calls to `\ExplSyntaxOff` are also not wasting time. Applying `\ExplSyntaxOn` alters the definition of `\ExplSyntaxOff` and so in package mode this function should not be used until after the end of the loading process!

```

147 \protected \def \ExplSyntaxOn
148   {
149     \bool_if:NF \l__kernel_expl_bool
150     {
151       \cs_set_protected:Npx \ExplSyntaxOff
152       {
153         \char_set_catcode:nn { 9 } { \char_value_catcode:n { 9 } }
154         \char_set_catcode:nn { 32 } { \char_value_catcode:n { 32 } }
155         \char_set_catcode:nn { 34 } { \char_value_catcode:n { 34 } }

```

```

156         \char_set_catcode:nn { 38 } { \char_value_catcode:n { 38 } }
157         \char_set_catcode:nn { 58 } { \char_value_catcode:n { 58 } }
158         \char_set_catcode:nn { 94 } { \char_value_catcode:n { 94 } }
159         \char_set_catcode:nn { 95 } { \char_value_catcode:n { 95 } }
160         \char_set_catcode:nn { 124 } { \char_value_catcode:n { 124 } }
161         \char_set_catcode:nn { 126 } { \char_value_catcode:n { 126 } }
162         \tex_endlinechar:D =
163         \tex_the:D \tex_endlinechar:D \scan_stop:
164         \bool_set_false:N \l__kernel_expl_bool
165         \cs_set_protected:Npn \ExplSyntaxOff { }
166     }
167 }
168 \char_set_catcode_ignore:n { 9 } % tab
169 \char_set_catcode_ignore:n { 32 } % space
170 \char_set_catcode_other:n { 34 } % double quote
171 \char_set_catcode_alignment:n { 38 } % ampersand
172 \char_set_catcode_letter:n { 58 } % colon
173 \char_set_catcode_math_superscript:n { 94 } % circumflex
174 \char_set_catcode_letter:n { 95 } % underscore
175 \char_set_catcode_other:n { 124 } % pipe
176 \char_set_catcode_space:n { 126 } % tilde
177 \tex_endlinechar:D = 32 \scan_stop:
178 \bool_set_true:N \l__kernel_expl_bool
179 }

```

(End definition for `\ExplSyntaxOn`. This function is documented on page 7.)

```

180 \endpackage

```

2 l3names implementation

```

181 \begin{package & tex}

```

The prefix here is `kernel`. A few places need `@@` to be left as is; this is obtained as `@@@`.

```

182 \@@=kernel

```

The code here simply renames all of the primitives to new, internal, names.

The `\let` primitive is renamed by hand first as it is essential for the entire process to follow. This also uses `\global`, as that way we avoid leaving an unneeded `csname` in the hash table.

```

183 \let \tex_global:D \global
184 \let \tex_let:D \let

```

Everything is inside a (rather long) group, which keeps `__kernel_primitive:NN` trapped.

```

185 \begin{group

```

`__kernel_primitive:NN` A temporary function to actually do the renaming.

```

186 \long \def \__kernel_primitive:NN #1#2
187 { \tex_global:D \tex_let:D #2 #1 }

```

(End definition for `__kernel_primitive:NN`.)

To allow extracting “just the names”, a bit of DocStrip fiddling.

```

188 \end{package & tex}
189 \begin{names | tex}
190 \begin{names | package}

```

In the current incarnation of this package, all TeX primitives are given a new name of the form `\tex_oldname:D`. But first three special cases which have symbolic original names. These are given modified new names, so that they may be entered without catcode tricks.

```

191 \__kernel_primitive:NN \tex_space:D
192 \__kernel_primitive:NN \tex_italiccorrection:D
193 \__kernel_primitive:NN \tex_hyphen:D

```

Now all the other primitives.

```

194 \__kernel_primitive:NN \tex_above:D
195 \__kernel_primitive:NN \tex_abovedisplayshortskip:D
196 \__kernel_primitive:NN \tex_abovedisplayskip:D
197 \__kernel_primitive:NN \tex_abovewithdelims:D
198 \__kernel_primitive:NN \tex_accent:D
199 \__kernel_primitive:NN \tex_adjdemerits:D
200 \__kernel_primitive:NN \tex_advance:D
201 \__kernel_primitive:NN \tex_afterassignment:D
202 \__kernel_primitive:NN \tex_aftergroup:D
203 \__kernel_primitive:NN \tex_atop:D
204 \__kernel_primitive:NN \tex_atopwithdelims:D
205 \__kernel_primitive:NN \tex_badness:D
206 \__kernel_primitive:NN \tex_baselineskip:D
207 \__kernel_primitive:NN \tex_batchmode:D
208 \__kernel_primitive:NN \tex_begingroup:D
209 \__kernel_primitive:NN \tex_belowdisplayshortskip:D
210 \__kernel_primitive:NN \tex_belowdisplayskip:D
211 \__kernel_primitive:NN \tex_binoppenalty:D
212 \__kernel_primitive:NN \tex_botmark:D
213 \__kernel_primitive:NN \tex_box:D
214 \__kernel_primitive:NN \tex_boxmaxdepth:D
215 \__kernel_primitive:NN \tex_brokenpenalty:D
216 \__kernel_primitive:NN \tex_catcode:D
217 \__kernel_primitive:NN \tex_char:D
218 \__kernel_primitive:NN \tex_chardef:D
219 \__kernel_primitive:NN \tex_cleaders:D
220 \__kernel_primitive:NN \tex_closein:D
221 \__kernel_primitive:NN \tex_closeout:D
222 \__kernel_primitive:NN \tex_clubpenalty:D
223 \__kernel_primitive:NN \tex_copy:D
224 \__kernel_primitive:NN \tex_count:D
225 \__kernel_primitive:NN \tex_countdef:D
226 \__kernel_primitive:NN \tex_cr:D
227 \__kernel_primitive:NN \tex_crcrcr:D
228 \__kernel_primitive:NN \tex_csname:D
229 \__kernel_primitive:NN \tex_day:D
230 \__kernel_primitive:NN \tex_deadcycles:D
231 \__kernel_primitive:NN \tex_def:D
232 \__kernel_primitive:NN \tex_defaultthyphenchar:D
233 \__kernel_primitive:NN \tex_defaultskewchar:D
234 \__kernel_primitive:NN \tex_delcode:D
235 \__kernel_primitive:NN \tex_delimiter:D
236 \__kernel_primitive:NN \tex_delimiterfactor:D
237 \__kernel_primitive:NN \tex_delimitershortfall:D
238 \__kernel_primitive:NN \tex_dimen:D

```

239	<code>__kernel_primitive:NN \dimendef</code>	<code>\tex_dimendef:D</code>
240	<code>__kernel_primitive:NN \discretionary</code>	<code>\tex_discretionary:D</code>
241	<code>__kernel_primitive:NN \displayindent</code>	<code>\tex_displayindent:D</code>
242	<code>__kernel_primitive:NN \displaylimits</code>	<code>\tex_displaylimits:D</code>
243	<code>__kernel_primitive:NN \displaystyle</code>	<code>\tex_displaystyle:D</code>
244	<code>__kernel_primitive:NN \displaywidowpenalty</code>	<code>\tex_displaywidowpenalty:D</code>
245	<code>__kernel_primitive:NN \displaywidth</code>	<code>\tex_displaywidth:D</code>
246	<code>__kernel_primitive:NN \divide</code>	<code>\tex_divide:D</code>
247	<code>__kernel_primitive:NN \doublehyphendemerits</code>	<code>\tex_doublehyphendemerits:D</code>
248	<code>__kernel_primitive:NN \dp</code>	<code>\tex_dp:D</code>
249	<code>__kernel_primitive:NN \dump</code>	<code>\tex_dump:D</code>
250	<code>__kernel_primitive:NN \edef</code>	<code>\tex_edef:D</code>
251	<code>__kernel_primitive:NN \else</code>	<code>\tex_else:D</code>
252	<code>__kernel_primitive:NN \emergencystretch</code>	<code>\tex_emergencystretch:D</code>
253	<code>__kernel_primitive:NN \end</code>	<code>\tex_end:D</code>
254	<code>__kernel_primitive:NN \endcsname</code>	<code>\tex_endcsname:D</code>
255	<code>__kernel_primitive:NN \endgroup</code>	<code>\tex_endgroup:D</code>
256	<code>__kernel_primitive:NN \endinput</code>	<code>\tex_endinput:D</code>
257	<code>__kernel_primitive:NN \endlinechar</code>	<code>\tex_endlinechar:D</code>
258	<code>__kernel_primitive:NN \eqno</code>	<code>\tex_eqno:D</code>
259	<code>__kernel_primitive:NN \errhelp</code>	<code>\tex_errhelp:D</code>
260	<code>__kernel_primitive:NN \errmessage</code>	<code>\tex_errmessage:D</code>
261	<code>__kernel_primitive:NN \errorcontextlines</code>	<code>\tex_errorcontextlines:D</code>
262	<code>__kernel_primitive:NN \errorstopmode</code>	<code>\tex_errorstopmode:D</code>
263	<code>__kernel_primitive:NN \escapechar</code>	<code>\tex_escapechar:D</code>
264	<code>__kernel_primitive:NN \everycr</code>	<code>\tex_everycr:D</code>
265	<code>__kernel_primitive:NN \everydisplay</code>	<code>\tex_everydisplay:D</code>
266	<code>__kernel_primitive:NN \everyhbox</code>	<code>\tex_everyhbox:D</code>
267	<code>__kernel_primitive:NN \everyjob</code>	<code>\tex_everyjob:D</code>
268	<code>__kernel_primitive:NN \everymath</code>	<code>\tex_everymath:D</code>
269	<code>__kernel_primitive:NN \everypar</code>	<code>\tex_everypar:D</code>
270	<code>__kernel_primitive:NN \everyvbox</code>	<code>\tex_everyvbox:D</code>
271	<code>__kernel_primitive:NN \exhyphenpenalty</code>	<code>\tex_exhyphenpenalty:D</code>
272	<code>__kernel_primitive:NN \expandafter</code>	<code>\tex_expandafter:D</code>
273	<code>__kernel_primitive:NN \fam</code>	<code>\tex_fam:D</code>
274	<code>__kernel_primitive:NN \fi</code>	<code>\tex_fi:D</code>
275	<code>__kernel_primitive:NN \finalhyphendemerits</code>	<code>\tex_finalhyphendemerits:D</code>
276	<code>__kernel_primitive:NN \firstmark</code>	<code>\tex_firstmark:D</code>
277	<code>__kernel_primitive:NN \floatingpenalty</code>	<code>\tex_floatingpenalty:D</code>
278	<code>__kernel_primitive:NN \font</code>	<code>\tex_font:D</code>
279	<code>__kernel_primitive:NN \fontdimen</code>	<code>\tex_fontdimen:D</code>
280	<code>__kernel_primitive:NN \fontname</code>	<code>\tex_fontname:D</code>
281	<code>__kernel_primitive:NN \futurelet</code>	<code>\tex_futurelet:D</code>
282	<code>__kernel_primitive:NN \gdef</code>	<code>\tex_gdef:D</code>
283	<code>__kernel_primitive:NN \global</code>	<code>\tex_global:D</code>
284	<code>__kernel_primitive:NN \globaldefs</code>	<code>\tex_globaldefs:D</code>
285	<code>__kernel_primitive:NN \halign</code>	<code>\tex_halign:D</code>
286	<code>__kernel_primitive:NN \hangafter</code>	<code>\tex_hangafter:D</code>
287	<code>__kernel_primitive:NN \hangindent</code>	<code>\tex_hangindent:D</code>
288	<code>__kernel_primitive:NN \hbadness</code>	<code>\tex_hbadness:D</code>
289	<code>__kernel_primitive:NN \hbox</code>	<code>\tex_hbox:D</code>
290	<code>__kernel_primitive:NN \hfil</code>	<code>\tex_hfil:D</code>
291	<code>__kernel_primitive:NN \hfill</code>	<code>\tex_hfill:D</code>
292	<code>__kernel_primitive:NN \hfilneg</code>	<code>\tex_hfilneg:D</code>

293	<code>__kernel_primitive:NN \hfuzz</code>	<code>\tex_hfuzz:D</code>
294	<code>__kernel_primitive:NN \hoffset</code>	<code>\tex_hoffset:D</code>
295	<code>__kernel_primitive:NN \holdinginserts</code>	<code>\tex_holdinginserts:D</code>
296	<code>__kernel_primitive:NN \hrule</code>	<code>\tex_hrule:D</code>
297	<code>__kernel_primitive:NN \hsize</code>	<code>\tex_hsize:D</code>
298	<code>__kernel_primitive:NN \hskip</code>	<code>\tex_hskip:D</code>
299	<code>__kernel_primitive:NN \hss</code>	<code>\tex_hss:D</code>
300	<code>__kernel_primitive:NN \ht</code>	<code>\tex_ht:D</code>
301	<code>__kernel_primitive:NN \hyphenation</code>	<code>\tex_hyphenation:D</code>
302	<code>__kernel_primitive:NN \hyphenchar</code>	<code>\tex_hyphenchar:D</code>
303	<code>__kernel_primitive:NN \hyphenpenalty</code>	<code>\tex_hyphenpenalty:D</code>
304	<code>__kernel_primitive:NN \if</code>	<code>\tex_if:D</code>
305	<code>__kernel_primitive:NN \ifcase</code>	<code>\tex_ifcase:D</code>
306	<code>__kernel_primitive:NN \ifcat</code>	<code>\tex_ifcat:D</code>
307	<code>__kernel_primitive:NN \ifdim</code>	<code>\tex_ifdim:D</code>
308	<code>__kernel_primitive:NN \ifeof</code>	<code>\tex_ifeof:D</code>
309	<code>__kernel_primitive:NN \iffalse</code>	<code>\tex_iffalse:D</code>
310	<code>__kernel_primitive:NN \ifhbox</code>	<code>\tex_ifhbox:D</code>
311	<code>__kernel_primitive:NN \ifhmode</code>	<code>\tex_ifhmode:D</code>
312	<code>__kernel_primitive:NN \ifinner</code>	<code>\tex_ifinner:D</code>
313	<code>__kernel_primitive:NN \ifmmode</code>	<code>\tex_ifmmode:D</code>
314	<code>__kernel_primitive:NN \ifnum</code>	<code>\tex_ifnum:D</code>
315	<code>__kernel_primitive:NN \ifodd</code>	<code>\tex_ifodd:D</code>
316	<code>__kernel_primitive:NN \iftrue</code>	<code>\tex_iftrue:D</code>
317	<code>__kernel_primitive:NN \ifvbox</code>	<code>\tex_ifvbox:D</code>
318	<code>__kernel_primitive:NN \ifvmode</code>	<code>\tex_ifvmode:D</code>
319	<code>__kernel_primitive:NN \ifvoid</code>	<code>\tex_ifvoid:D</code>
320	<code>__kernel_primitive:NN \ifx</code>	<code>\tex_ifx:D</code>
321	<code>__kernel_primitive:NN \ignorespaces</code>	<code>\tex_ignorespaces:D</code>
322	<code>__kernel_primitive:NN \immediate</code>	<code>\tex_immediate:D</code>
323	<code>__kernel_primitive:NN \indent</code>	<code>\tex_indent:D</code>
324	<code>__kernel_primitive:NN \input</code>	<code>\tex_input:D</code>
325	<code>__kernel_primitive:NN \inputlineno</code>	<code>\tex_inputlineno:D</code>
326	<code>__kernel_primitive:NN \insert</code>	<code>\tex_insert:D</code>
327	<code>__kernel_primitive:NN \insertpenalties</code>	<code>\tex_insertpenalties:D</code>
328	<code>__kernel_primitive:NN \interlinepenalty</code>	<code>\tex_interlinepenalty:D</code>
329	<code>__kernel_primitive:NN \jobname</code>	<code>\tex_jobname:D</code>
330	<code>__kernel_primitive:NN \kern</code>	<code>\tex_kern:D</code>
331	<code>__kernel_primitive:NN \language</code>	<code>\tex_language:D</code>
332	<code>__kernel_primitive:NN \lastbox</code>	<code>\tex_lastbox:D</code>
333	<code>__kernel_primitive:NN \lastkern</code>	<code>\tex_lastkern:D</code>
334	<code>__kernel_primitive:NN \lastpenalty</code>	<code>\tex_lastpenalty:D</code>
335	<code>__kernel_primitive:NN \lastskip</code>	<code>\tex_lastskip:D</code>
336	<code>__kernel_primitive:NN \lccode</code>	<code>\tex_lccode:D</code>
337	<code>__kernel_primitive:NN \leaders</code>	<code>\tex_leaders:D</code>
338	<code>__kernel_primitive:NN \left</code>	<code>\tex_left:D</code>
339	<code>__kernel_primitive:NN \lefthyphenmin</code>	<code>\tex_lefthyphenmin:D</code>
340	<code>__kernel_primitive:NN \leftskip</code>	<code>\tex_leftskip:D</code>
341	<code>__kernel_primitive:NN \leqno</code>	<code>\tex_leqno:D</code>
342	<code>__kernel_primitive:NN \let</code>	<code>\tex_let:D</code>
343	<code>__kernel_primitive:NN \limits</code>	<code>\tex_limits:D</code>
344	<code>__kernel_primitive:NN \linepenalty</code>	<code>\tex_linepenalty:D</code>
345	<code>__kernel_primitive:NN \lineskip</code>	<code>\tex_lineskip:D</code>
346	<code>__kernel_primitive:NN \lineskiplimit</code>	<code>\tex_lineskiplimit:D</code>

347	<code>__kernel_primitive:NN \long</code>	<code>\tex_long:D</code>
348	<code>__kernel_primitive:NN \looseness</code>	<code>\tex_looseness:D</code>
349	<code>__kernel_primitive:NN \lower</code>	<code>\tex_lower:D</code>
350	<code>__kernel_primitive:NN \lowercase</code>	<code>\tex_lowercase:D</code>
351	<code>__kernel_primitive:NN \mag</code>	<code>\tex_mag:D</code>
352	<code>__kernel_primitive:NN \mark</code>	<code>\tex_mark:D</code>
353	<code>__kernel_primitive:NN \mathaccent</code>	<code>\tex_mathaccent:D</code>
354	<code>__kernel_primitive:NN \mathbin</code>	<code>\tex_mathbin:D</code>
355	<code>__kernel_primitive:NN \mathchar</code>	<code>\tex_mathchar:D</code>
356	<code>__kernel_primitive:NN \mathchardef</code>	<code>\tex_mathchardef:D</code>
357	<code>__kernel_primitive:NN \mathchoice</code>	<code>\tex_mathchoice:D</code>
358	<code>__kernel_primitive:NN \mathclose</code>	<code>\tex_mathclose:D</code>
359	<code>__kernel_primitive:NN \mathcode</code>	<code>\tex_mathcode:D</code>
360	<code>__kernel_primitive:NN \mathinner</code>	<code>\tex_mathinner:D</code>
361	<code>__kernel_primitive:NN \mathop</code>	<code>\tex_mathop:D</code>
362	<code>__kernel_primitive:NN \mathopen</code>	<code>\tex_mathopen:D</code>
363	<code>__kernel_primitive:NN \mathord</code>	<code>\tex_mathord:D</code>
364	<code>__kernel_primitive:NN \mathpunct</code>	<code>\tex_mathpunct:D</code>
365	<code>__kernel_primitive:NN \mathrel</code>	<code>\tex_mathrel:D</code>
366	<code>__kernel_primitive:NN \mathsurround</code>	<code>\tex_mathsurround:D</code>
367	<code>__kernel_primitive:NN \maxdeadcycles</code>	<code>\tex_maxdeadcycles:D</code>
368	<code>__kernel_primitive:NN \maxdepth</code>	<code>\tex_maxdepth:D</code>
369	<code>__kernel_primitive:NN \meaning</code>	<code>\tex_meaning:D</code>
370	<code>__kernel_primitive:NN \medmuskip</code>	<code>\tex_medmuskip:D</code>
371	<code>__kernel_primitive:NN \message</code>	<code>\tex_message:D</code>
372	<code>__kernel_primitive:NN \mkern</code>	<code>\tex_mkern:D</code>
373	<code>__kernel_primitive:NN \month</code>	<code>\tex_month:D</code>
374	<code>__kernel_primitive:NN \moveleft</code>	<code>\tex_moveleft:D</code>
375	<code>__kernel_primitive:NN \moveright</code>	<code>\tex_moveright:D</code>
376	<code>__kernel_primitive:NN \mskip</code>	<code>\tex_mskip:D</code>
377	<code>__kernel_primitive:NN \multiply</code>	<code>\tex_multiply:D</code>
378	<code>__kernel_primitive:NN \muskip</code>	<code>\tex_muskip:D</code>
379	<code>__kernel_primitive:NN \muskipdef</code>	<code>\tex_muskipdef:D</code>
380	<code>__kernel_primitive:NN \newlinechar</code>	<code>\tex_newlinechar:D</code>
381	<code>__kernel_primitive:NN \noalign</code>	<code>\tex_noalign:D</code>
382	<code>__kernel_primitive:NN \noboundary</code>	<code>\tex_noboundary:D</code>
383	<code>__kernel_primitive:NN \noexpand</code>	<code>\tex_noexpand:D</code>
384	<code>__kernel_primitive:NN \noindent</code>	<code>\tex_noindent:D</code>
385	<code>__kernel_primitive:NN \nolimits</code>	<code>\tex_nolimits:D</code>
386	<code>__kernel_primitive:NN \nonscript</code>	<code>\tex_nonscript:D</code>
387	<code>__kernel_primitive:NN \nonstopmode</code>	<code>\tex_nonstopmode:D</code>
388	<code>__kernel_primitive:NN \nulldelimiterspace</code>	<code>\tex_nulldelimiterspace:D</code>
389	<code>__kernel_primitive:NN \nullfont</code>	<code>\tex_nullfont:D</code>
390	<code>__kernel_primitive:NN \number</code>	<code>\tex_number:D</code>
391	<code>__kernel_primitive:NN \omit</code>	<code>\tex_omit:D</code>
392	<code>__kernel_primitive:NN \openin</code>	<code>\tex_openin:D</code>
393	<code>__kernel_primitive:NN \openout</code>	<code>\tex_openout:D</code>
394	<code>__kernel_primitive:NN \or</code>	<code>\tex_or:D</code>
395	<code>__kernel_primitive:NN \outer</code>	<code>\tex_outer:D</code>
396	<code>__kernel_primitive:NN \output</code>	<code>\tex_output:D</code>
397	<code>__kernel_primitive:NN \outputpenalty</code>	<code>\tex_outputpenalty:D</code>
398	<code>__kernel_primitive:NN \over</code>	<code>\tex_over:D</code>
399	<code>__kernel_primitive:NN \overfullrule</code>	<code>\tex_overfullrule:D</code>
400	<code>__kernel_primitive:NN \overline</code>	<code>\tex_overline:D</code>

401	<code>__kernel_primitive:NN \overwithdelims</code>	<code>\tex_overwithdelims:D</code>
402	<code>__kernel_primitive:NN \pagedepth</code>	<code>\tex_pagedepth:D</code>
403	<code>__kernel_primitive:NN \pagefilllstretch</code>	<code>\tex_pagefilllstretch:D</code>
404	<code>__kernel_primitive:NN \pagefillstretch</code>	<code>\tex_pagefillstretch:D</code>
405	<code>__kernel_primitive:NN \pagefilstretch</code>	<code>\tex_pagefilstretch:D</code>
406	<code>__kernel_primitive:NN \pagegoal</code>	<code>\tex_pagegoal:D</code>
407	<code>__kernel_primitive:NN \pageshrink</code>	<code>\tex_pageshrink:D</code>
408	<code>__kernel_primitive:NN \pagestretch</code>	<code>\tex_pagestretch:D</code>
409	<code>__kernel_primitive:NN \pagetotal</code>	<code>\tex_pagetotal:D</code>
410	<code>__kernel_primitive:NN \par</code>	<code>\tex_par:D</code>
411	<code>__kernel_primitive:NN \parfillskip</code>	<code>\tex_parfillskip:D</code>
412	<code>__kernel_primitive:NN \parindent</code>	<code>\tex_parindent:D</code>
413	<code>__kernel_primitive:NN \parshape</code>	<code>\tex_parshape:D</code>
414	<code>__kernel_primitive:NN \parskip</code>	<code>\tex_parskip:D</code>
415	<code>__kernel_primitive:NN \patterns</code>	<code>\tex_patterns:D</code>
416	<code>__kernel_primitive:NN \pausing</code>	<code>\tex_pausing:D</code>
417	<code>__kernel_primitive:NN \penalty</code>	<code>\tex_penalty:D</code>
418	<code>__kernel_primitive:NN \postdisplaypenalty</code>	<code>\tex_postdisplaypenalty:D</code>
419	<code>__kernel_primitive:NN \predisdisplaypenalty</code>	<code>\tex_predisdisplaypenalty:D</code>
420	<code>__kernel_primitive:NN \predisplaysize</code>	<code>\tex_predisplaysize:D</code>
421	<code>__kernel_primitive:NN \pretolerance</code>	<code>\tex_pretolerance:D</code>
422	<code>__kernel_primitive:NN \prevdepth</code>	<code>\tex_prevdepth:D</code>
423	<code>__kernel_primitive:NN \prevgraf</code>	<code>\tex_prevgraf:D</code>
424	<code>__kernel_primitive:NN \radical</code>	<code>\tex_radical:D</code>
425	<code>__kernel_primitive:NN \raise</code>	<code>\tex_raise:D</code>
426	<code>__kernel_primitive:NN \read</code>	<code>\tex_read:D</code>
427	<code>__kernel_primitive:NN \relax</code>	<code>\tex_relax:D</code>
428	<code>__kernel_primitive:NN \relpenalty</code>	<code>\tex_relpenalty:D</code>
429	<code>__kernel_primitive:NN \right</code>	<code>\tex_right:D</code>
430	<code>__kernel_primitive:NN \righthyphenmin</code>	<code>\tex_righthyphenmin:D</code>
431	<code>__kernel_primitive:NN \rightskip</code>	<code>\tex_rightskip:D</code>
432	<code>__kernel_primitive:NN \romannumeral</code>	<code>\tex_romannumeral:D</code>
433	<code>__kernel_primitive:NN \scriptfont</code>	<code>\tex_scriptfont:D</code>
434	<code>__kernel_primitive:NN \scriptscriptfont</code>	<code>\tex_scriptscriptfont:D</code>
435	<code>__kernel_primitive:NN \scriptscriptstyle</code>	<code>\tex_scriptscriptstyle:D</code>
436	<code>__kernel_primitive:NN \scriptspace</code>	<code>\tex_scriptspace:D</code>
437	<code>__kernel_primitive:NN \scriptstyle</code>	<code>\tex_scriptstyle:D</code>
438	<code>__kernel_primitive:NN \scrollmode</code>	<code>\tex_scrollmode:D</code>
439	<code>__kernel_primitive:NN \setbox</code>	<code>\tex_setbox:D</code>
440	<code>__kernel_primitive:NN \setlanguage</code>	<code>\tex_setlanguage:D</code>
441	<code>__kernel_primitive:NN \sfcode</code>	<code>\tex_sfcode:D</code>
442	<code>__kernel_primitive:NN \shipout</code>	<code>\tex_shipout:D</code>
443	<code>__kernel_primitive:NN \show</code>	<code>\tex_show:D</code>
444	<code>__kernel_primitive:NN \showbox</code>	<code>\tex_showbox:D</code>
445	<code>__kernel_primitive:NN \showboxbreadth</code>	<code>\tex_showboxbreadth:D</code>
446	<code>__kernel_primitive:NN \showboxdepth</code>	<code>\tex_showboxdepth:D</code>
447	<code>__kernel_primitive:NN \showlists</code>	<code>\tex_showlists:D</code>
448	<code>__kernel_primitive:NN \showthe</code>	<code>\tex_showthe:D</code>
449	<code>__kernel_primitive:NN \skewchar</code>	<code>\tex_skewchar:D</code>
450	<code>__kernel_primitive:NN \skip</code>	<code>\tex_skip:D</code>
451	<code>__kernel_primitive:NN \skipdef</code>	<code>\tex_skipdef:D</code>
452	<code>__kernel_primitive:NN \spacefactor</code>	<code>\tex_spacefactor:D</code>
453	<code>__kernel_primitive:NN \spaceskip</code>	<code>\tex_spaceskip:D</code>
454	<code>__kernel_primitive:NN \span</code>	<code>\tex_span:D</code>

455	_kernel_primitive:NN	\special	\tex_special:D
456	_kernel_primitive:NN	\splitbotmark	\tex_splitbotmark:D
457	_kernel_primitive:NN	\splitfirstmark	\tex_splitfirstmark:D
458	_kernel_primitive:NN	\splitmaxdepth	\tex_splitmaxdepth:D
459	_kernel_primitive:NN	\splittopskip	\tex_splittopskip:D
460	_kernel_primitive:NN	\string	\tex_string:D
461	_kernel_primitive:NN	\tabskip	\tex_tabskip:D
462	_kernel_primitive:NN	\textfont	\tex_textfont:D
463	_kernel_primitive:NN	\textstyle	\tex_textstyle:D
464	_kernel_primitive:NN	\the	\tex_the:D
465	_kernel_primitive:NN	\thickmuskip	\tex_thickmuskip:D
466	_kernel_primitive:NN	\thinmuskip	\tex_thinmuskip:D
467	_kernel_primitive:NN	\time	\tex_time:D
468	_kernel_primitive:NN	\toks	\tex_toks:D
469	_kernel_primitive:NN	\toksdef	\tex_toksdef:D
470	_kernel_primitive:NN	\tolerance	\tex_tolerance:D
471	_kernel_primitive:NN	\topmark	\tex_topmark:D
472	_kernel_primitive:NN	\topskip	\tex_topskip:D
473	_kernel_primitive:NN	\tracingcommands	\tex_tracingcommands:D
474	_kernel_primitive:NN	\tracinglostchars	\tex_tracinglostchars:D
475	_kernel_primitive:NN	\tracingmacros	\tex_tracingmacros:D
476	_kernel_primitive:NN	\tracingonline	\tex_tracingonline:D
477	_kernel_primitive:NN	\tracingoutput	\tex_tracingoutput:D
478	_kernel_primitive:NN	\tracingpages	\tex_tracingpages:D
479	_kernel_primitive:NN	\tracingparagraphs	\tex_tracingparagraphs:D
480	_kernel_primitive:NN	\tracingrestores	\tex_tracingrestores:D
481	_kernel_primitive:NN	\tracingstats	\tex_tracingstats:D
482	_kernel_primitive:NN	\uccode	\tex_uccode:D
483	_kernel_primitive:NN	\uchyph	\tex_uchyph:D
484	_kernel_primitive:NN	\underline	\tex_underline:D
485	_kernel_primitive:NN	\unhbox	\tex_unhbox:D
486	_kernel_primitive:NN	\unhcopy	\tex_unhcopy:D
487	_kernel_primitive:NN	\unkern	\tex_unkern:D
488	_kernel_primitive:NN	\unpenalty	\tex_unpenalty:D
489	_kernel_primitive:NN	\unskip	\tex_unskip:D
490	_kernel_primitive:NN	\unvbox	\tex_unvbox:D
491	_kernel_primitive:NN	\unvcopy	\tex_unvcopy:D
492	_kernel_primitive:NN	\uppercase	\tex_uppercase:D
493	_kernel_primitive:NN	\vadjust	\tex_vadjust:D
494	_kernel_primitive:NN	\valign	\tex_valign:D
495	_kernel_primitive:NN	\vbadness	\tex_vbadness:D
496	_kernel_primitive:NN	\vbox	\tex_vbox:D
497	_kernel_primitive:NN	\vcenter	\tex_vcenter:D
498	_kernel_primitive:NN	\vfil	\tex_vfil:D
499	_kernel_primitive:NN	\vfill	\tex_vfill:D
500	_kernel_primitive:NN	\vfилneg	\tex_vfилneg:D
501	_kernel_primitive:NN	\vfuzz	\tex_vfuzz:D
502	_kernel_primitive:NN	\voffset	\tex_voffset:D
503	_kernel_primitive:NN	\vrule	\tex_vrule:D
504	_kernel_primitive:NN	\vsize	\tex_vsize:D
505	_kernel_primitive:NN	\vskip	\tex_vskip:D
506	_kernel_primitive:NN	\vsplit	\tex_vsplit:D
507	_kernel_primitive:NN	\vss	\tex_vss:D
508	_kernel_primitive:NN	\vtop	\tex_vtop:D

509	<code>_kernel_primitive:NN \wd</code>	<code>\tex_wd:D</code>
510	<code>_kernel_primitive:NN \widowpenalty</code>	<code>\tex_widowpenalty:D</code>
511	<code>_kernel_primitive:NN \write</code>	<code>\tex_write:D</code>
512	<code>_kernel_primitive:NN \xdef</code>	<code>\tex_xdef:D</code>
513	<code>_kernel_primitive:NN \xleaders</code>	<code>\tex_xleaders:D</code>
514	<code>_kernel_primitive:NN \xspaceskip</code>	<code>\tex_xspaceskip:D</code>
515	<code>_kernel_primitive:NN \year</code>	<code>\tex_year:D</code>

Primitives introduced by ε -T_EX.

516	<code>_kernel_primitive:NN \beginL</code>	<code>\tex_beginL:D</code>
517	<code>_kernel_primitive:NN \beginR</code>	<code>\tex_beginR:D</code>
518	<code>_kernel_primitive:NN \botmarks</code>	<code>\tex_botmarks:D</code>
519	<code>_kernel_primitive:NN \clubpenalties</code>	<code>\tex_clubpenalties:D</code>
520	<code>_kernel_primitive:NN \currentgrouplevel</code>	<code>\tex_currentgrouplevel:D</code>
521	<code>_kernel_primitive:NN \currentgrouptype</code>	<code>\tex_currentgrouptype:D</code>
522	<code>_kernel_primitive:NN \currentifbranch</code>	<code>\tex_currentifbranch:D</code>
523	<code>_kernel_primitive:NN \currentiflevel</code>	<code>\tex_currentiflevel:D</code>
524	<code>_kernel_primitive:NN \currentifttype</code>	<code>\tex_currentifttype:D</code>
525	<code>_kernel_primitive:NN \detokenize</code>	<code>\tex_detokenize:D</code>
526	<code>_kernel_primitive:NN \dimexpr</code>	<code>\tex_dimexpr:D</code>
527	<code>_kernel_primitive:NN \displaywidowpenalties</code>	<code>\tex_displaywidowpenalties:D</code>
528	<code>_kernel_primitive:NN \endL</code>	<code>\tex_endL:D</code>
529	<code>_kernel_primitive:NN \endR</code>	<code>\tex_endR:D</code>
530	<code>_kernel_primitive:NN \eTeXrevision</code>	<code>\tex_eTeXrevision:D</code>
531	<code>_kernel_primitive:NN \eTeXversion</code>	<code>\tex_eTeXversion:D</code>
532	<code>_kernel_primitive:NN \everyeof</code>	<code>\tex_everyeof:D</code>
533	<code>_kernel_primitive:NN \firstmarks</code>	<code>\tex_firstmarks:D</code>
534	<code>_kernel_primitive:NN \fontchardp</code>	<code>\tex_fontchardp:D</code>
535	<code>_kernel_primitive:NN \fontcharht</code>	<code>\tex_fontcharht:D</code>
536	<code>_kernel_primitive:NN \fontcharic</code>	<code>\tex_fontcharic:D</code>
537	<code>_kernel_primitive:NN \fontcharwd</code>	<code>\tex_fontcharwd:D</code>
538	<code>_kernel_primitive:NN \glueexpr</code>	<code>\tex_glueexpr:D</code>
539	<code>_kernel_primitive:NN \glueshrink</code>	<code>\tex_glueshrink:D</code>
540	<code>_kernel_primitive:NN \glueshrinkorder</code>	<code>\tex_glueshrinkorder:D</code>
541	<code>_kernel_primitive:NN \gluestretch</code>	<code>\tex_gluestretch:D</code>
542	<code>_kernel_primitive:NN \gluestretchorder</code>	<code>\tex_gluestretchorder:D</code>
543	<code>_kernel_primitive:NN \gluetomu</code>	<code>\tex_gluetomu:D</code>
544	<code>_kernel_primitive:NN \ifcsname</code>	<code>\tex_ifcsname:D</code>
545	<code>_kernel_primitive:NN \ifdefined</code>	<code>\tex_ifdefined:D</code>
546	<code>_kernel_primitive:NN \iffontchar</code>	<code>\tex_iffontchar:D</code>
547	<code>_kernel_primitive:NN \interactionmode</code>	<code>\tex_interactionmode:D</code>
548	<code>_kernel_primitive:NN \interlinepenalties</code>	<code>\tex_interlinepenalties:D</code>
549	<code>_kernel_primitive:NN \lastlinefit</code>	<code>\tex_lastlinefit:D</code>
550	<code>_kernel_primitive:NN \lastnodetype</code>	<code>\tex_lastnodetype:D</code>
551	<code>_kernel_primitive:NN \marks</code>	<code>\tex_marks:D</code>
552	<code>_kernel_primitive:NN \middle</code>	<code>\tex_middle:D</code>
553	<code>_kernel_primitive:NN \muexpr</code>	<code>\tex_muexpr:D</code>
554	<code>_kernel_primitive:NN \mutoglu</code>	<code>\tex_mutoglu:D</code>
555	<code>_kernel_primitive:NN \numexpr</code>	<code>\tex_numexpr:D</code>
556	<code>_kernel_primitive:NN \pagediscards</code>	<code>\tex_pagediscards:D</code>
557	<code>_kernel_primitive:NN \parshapedimen</code>	<code>\tex_parshapedimen:D</code>
558	<code>_kernel_primitive:NN \parshapeindent</code>	<code>\tex_parshapeindent:D</code>
559	<code>_kernel_primitive:NN \parshapelength</code>	<code>\tex_parshapelength:D</code>
560	<code>_kernel_primitive:NN \predisplaydirection</code>	<code>\tex_predisplaydirection:D</code>
561	<code>_kernel_primitive:NN \protected</code>	<code>\tex_protected:D</code>

562	<code>__kernel_primitive:NN \readline</code>	<code>\tex_readline:D</code>
563	<code>__kernel_primitive:NN \savinghyphcodes</code>	<code>\tex_savinghyphcodes:D</code>
564	<code>__kernel_primitive:NN \savingvdiscards</code>	<code>\tex_savingvdiscards:D</code>
565	<code>__kernel_primitive:NN \scantokens</code>	<code>\tex_scantokens:D</code>
566	<code>__kernel_primitive:NN \showgroups</code>	<code>\tex_showgroups:D</code>
567	<code>__kernel_primitive:NN \showifs</code>	<code>\tex_showifs:D</code>
568	<code>__kernel_primitive:NN \showtokens</code>	<code>\tex_showtokens:D</code>
569	<code>__kernel_primitive:NN \splitbotmarks</code>	<code>\tex_splitbotmarks:D</code>
570	<code>__kernel_primitive:NN \splitdiscards</code>	<code>\tex_splitdiscards:D</code>
571	<code>__kernel_primitive:NN \splitfirstmarks</code>	<code>\tex_splitfirstmarks:D</code>
572	<code>__kernel_primitive:NN \TeXeTstate</code>	<code>\tex_TeXeTstate:D</code>
573	<code>__kernel_primitive:NN \topmarks</code>	<code>\tex_topmarks:D</code>
574	<code>__kernel_primitive:NN \tracingassigns</code>	<code>\tex_tracingassigns:D</code>
575	<code>__kernel_primitive:NN \tracinggroups</code>	<code>\tex_tracinggroups:D</code>
576	<code>__kernel_primitive:NN \tracingifs</code>	<code>\tex_tracingifs:D</code>
577	<code>__kernel_primitive:NN \tracingnesting</code>	<code>\tex_tracingnesting:D</code>
578	<code>__kernel_primitive:NN \tracingscantokens</code>	<code>\tex_tracingscantokens:D</code>
579	<code>__kernel_primitive:NN \unexpanded</code>	<code>\tex_unexpanded:D</code>
580	<code>__kernel_primitive:NN \unless</code>	<code>\tex_unless:D</code>
581	<code>__kernel_primitive:NN \widowpenalties</code>	<code>\tex_widowpenalties:D</code>

Post- ϵ -TeX primitives do not always end up with the same name in all engines, if indeed they are available cross-engine anyway. We therefore take the approach of preferring the shortest name that makes sense. First, we deal with the primitives introduced by pdfTeX which directly relate to PDF output: these are copied with the names unchanged.

582	<code>__kernel_primitive:NN \pdfannot</code>	<code>\tex_pdfannot:D</code>
583	<code>__kernel_primitive:NN \pdfcatalog</code>	<code>\tex_pdfcatalog:D</code>
584	<code>__kernel_primitive:NN \pdfcompresslevel</code>	<code>\tex_pdfcompresslevel:D</code>
585	<code>__kernel_primitive:NN \pdfcolorstack</code>	<code>\tex_pdfcolorstack:D</code>
586	<code>__kernel_primitive:NN \pdfcolorstackinit</code>	<code>\tex_pdfcolorstackinit:D</code>
587	<code>__kernel_primitive:NN \pdfcreationdate</code>	<code>\tex_pdfcreationdate:D</code>
588	<code>__kernel_primitive:NN \pdfdecimaldigits</code>	<code>\tex_pdfdecimaldigits:D</code>
589	<code>__kernel_primitive:NN \pdfdest</code>	<code>\tex_pdfdest:D</code>
590	<code>__kernel_primitive:NN \pdfdestmargin</code>	<code>\tex_pdfdestmargin:D</code>
591	<code>__kernel_primitive:NN \pdfendlink</code>	<code>\tex_pdfendlink:D</code>
592	<code>__kernel_primitive:NN \pdfendthread</code>	<code>\tex_pdfendthread:D</code>
593	<code>__kernel_primitive:NN \pdffontattr</code>	<code>\tex_pdffontattr:D</code>
594	<code>__kernel_primitive:NN \pdffontname</code>	<code>\tex_pdffontname:D</code>
595	<code>__kernel_primitive:NN \pdffontobjnum</code>	<code>\tex_pdffontobjnum:D</code>
596	<code>__kernel_primitive:NN \pdfgamma</code>	<code>\tex_pdfgamma:D</code>
597	<code>__kernel_primitive:NN \pdfimageapplygamma</code>	<code>\tex_pdfimageapplygamma:D</code>
598	<code>__kernel_primitive:NN \pdfimagegamma</code>	<code>\tex_pdfimagegamma:D</code>
599	<code>__kernel_primitive:NN \pdfgentounicode</code>	<code>\tex_pdfgentounicode:D</code>
600	<code>__kernel_primitive:NN \pdfglyptounicode</code>	<code>\tex_pdfglyptounicode:D</code>
601	<code>__kernel_primitive:NN \pdfhorigin</code>	<code>\tex_pdfhorigin:D</code>
602	<code>__kernel_primitive:NN \pdfimagehicolor</code>	<code>\tex_pdfimagehicolor:D</code>
603	<code>__kernel_primitive:NN \pdfimageresolution</code>	<code>\tex_pdfimageresolution:D</code>
604	<code>__kernel_primitive:NN \pdfincludechars</code>	<code>\tex_pdfincludechars:D</code>
605	<code>__kernel_primitive:NN \pdfinclusioncopyfonts</code>	<code>\tex_pdfinclusioncopyfonts:D</code>
606	<code>__kernel_primitive:NN \pdfinclusionerrorlevel</code>	
607	<code>\tex_pdfinclusionerrorlevel:D</code>	
608	<code>__kernel_primitive:NN \pdfinfo</code>	<code>\tex_pdfinfo:D</code>
609	<code>__kernel_primitive:NN \pdflastannot</code>	<code>\tex_pdflastannot:D</code>
610	<code>__kernel_primitive:NN \pdflastlink</code>	<code>\tex_pdflastlink:D</code>

611	<code>__kernel_primitive:NN \pdflastobj</code>	<code>\tex_pdflastobj:D</code>
612	<code>__kernel_primitive:NN \pdflastxform</code>	<code>\tex_pdflastxform:D</code>
613	<code>__kernel_primitive:NN \pdflastximage</code>	<code>\tex_pdflastximage:D</code>
614	<code>__kernel_primitive:NN \pdflastximagecolordepth</code>	
615	<code>\tex_pdflastximagecolordepth:D</code>	
616	<code>__kernel_primitive:NN \pdflastximagepages</code>	<code>\tex_pdflastximagepages:D</code>
617	<code>__kernel_primitive:NN \pdflinkmargin</code>	<code>\tex_pdflinkmargin:D</code>
618	<code>__kernel_primitive:NN \pdfliteral</code>	<code>\tex_pdfliteral:D</code>
619	<code>__kernel_primitive:NN \pdfmajorversion</code>	<code>\tex_pdfmajorversion:D</code>
620	<code>__kernel_primitive:NN \pdfminorversion</code>	<code>\tex_pdfminorversion:D</code>
621	<code>__kernel_primitive:NN \pdfnames</code>	<code>\tex_pdfnames:D</code>
622	<code>__kernel_primitive:NN \pdfobj</code>	<code>\tex_pdfobj:D</code>
623	<code>__kernel_primitive:NN \pdfobjcompresslevel</code>	<code>\tex_pdfobjcompresslevel:D</code>
624	<code>__kernel_primitive:NN \pdfoutline</code>	<code>\tex_pdfoutline:D</code>
625	<code>__kernel_primitive:NN \pdfoutput</code>	<code>\tex_pdfoutput:D</code>
626	<code>__kernel_primitive:NN \pdfpageattr</code>	<code>\tex_pdfpageattr:D</code>
627	<code>__kernel_primitive:NN \pdfpagesattr</code>	<code>\tex_pdfpagesattr:D</code>
628	<code>__kernel_primitive:NN \pdfpagebox</code>	<code>\tex_pdfpagebox:D</code>
629	<code>__kernel_primitive:NN \pdfpageref</code>	<code>\tex_pdfpageref:D</code>
630	<code>__kernel_primitive:NN \pdfpageresources</code>	<code>\tex_pdfpageresources:D</code>
631	<code>__kernel_primitive:NN \pdfpagesattr</code>	<code>\tex_pdfpagesattr:D</code>
632	<code>__kernel_primitive:NN \pdfrefobj</code>	<code>\tex_pdfrefobj:D</code>
633	<code>__kernel_primitive:NN \pdfrefxform</code>	<code>\tex_pdfrefxform:D</code>
634	<code>__kernel_primitive:NN \pdfrefximage</code>	<code>\tex_pdfrefximage:D</code>
635	<code>__kernel_primitive:NN \pdfrestore</code>	<code>\tex_pdfrestore:D</code>
636	<code>__kernel_primitive:NN \pdfretval</code>	<code>\tex_pdfretval:D</code>
637	<code>__kernel_primitive:NN \pdfsave</code>	<code>\tex_pdfsave:D</code>
638	<code>__kernel_primitive:NN \pdfsetmatrix</code>	<code>\tex_pdfsetmatrix:D</code>
639	<code>__kernel_primitive:NN \pdfstartlink</code>	<code>\tex_pdfstartlink:D</code>
640	<code>__kernel_primitive:NN \pdfstartthread</code>	<code>\tex_pdfstartthread:D</code>
641	<code>__kernel_primitive:NN \pdfsuppressptexinfo</code>	<code>\tex_pdfsuppressptexinfo:D</code>
642	<code>__kernel_primitive:NN \pdfthread</code>	<code>\tex_pdfthread:D</code>
643	<code>__kernel_primitive:NN \pdfthreadmargin</code>	<code>\tex_pdfthreadmargin:D</code>
644	<code>__kernel_primitive:NN \pdftrailer</code>	<code>\tex_pdftrailer:D</code>
645	<code>__kernel_primitive:NN \pdfuniquestring</code>	<code>\tex_pdfuniquestring:D</code>
646	<code>__kernel_primitive:NN \pdfvorigin</code>	<code>\tex_pdfvorigin:D</code>
647	<code>__kernel_primitive:NN \pdfxform</code>	<code>\tex_pdfxform:D</code>
648	<code>__kernel_primitive:NN \pdfxformname</code>	<code>\tex_pdfxformname:D</code>
649	<code>__kernel_primitive:NN \pdfximage</code>	<code>\tex_pdfximage:D</code>
650	<code>__kernel_primitive:NN \pdfximagebbox</code>	<code>\tex_pdfximagebbox:D</code>

These are not related to PDF output and either already appear in other engines without the `\pdf` prefix, or might reasonably do so at some future stage. We therefore drop the leading `pdf` here.

651	<code>__kernel_primitive:NN \ifpdfabsdim</code>	<code>\tex_ifabsdim:D</code>
652	<code>__kernel_primitive:NN \ifpdfabsnum</code>	<code>\tex_ifabsnum:D</code>
653	<code>__kernel_primitive:NN \ifpdfprimitive</code>	<code>\tex_ifprimitive:D</code>
654	<code>__kernel_primitive:NN \pdfadjustspacing</code>	<code>\tex_adjustspacing:D</code>
655	<code>__kernel_primitive:NN \pdfcopyfont</code>	<code>\tex_copyfont:D</code>
656	<code>__kernel_primitive:NN \pdfdraftmode</code>	<code>\tex_draftmode:D</code>
657	<code>__kernel_primitive:NN \pdfeachlinedepth</code>	<code>\tex_eachlinedepth:D</code>
658	<code>__kernel_primitive:NN \pdfeachlineheight</code>	<code>\tex_eachlineheight:D</code>
659	<code>__kernel_primitive:NN \pdfelapsedtime</code>	<code>\tex_elapsedtime:D</code>
660	<code>__kernel_primitive:NN \pdffirstlineheight</code>	<code>\tex_firstlineheight:D</code>

```

661 \__kernel_primitive:NN \pdffontexpand \tex_fontexpand:D
662 \__kernel_primitive:NN \pdffontsize \tex_fontsize:D
663 \__kernel_primitive:NN \pdfignoreddimen \tex_ignoreddimen:D
664 \__kernel_primitive:NN \pdfinsertht \tex_insertht:D
665 \__kernel_primitive:NN \pdflastlinedepth \tex_lastlinedepth:D
666 \__kernel_primitive:NN \pdflastxpos \tex_lastxpos:D
667 \__kernel_primitive:NN \pdflastypos \tex_lastypos:D
668 \__kernel_primitive:NN \pdfmapfile \tex_mapfile:D
669 \__kernel_primitive:NN \pdfmapline \tex_mapline:D
670 \__kernel_primitive:NN \pdfnoligatures \tex_noligatures:D
671 \__kernel_primitive:NN \pdfnormaldeviate \tex_normaldeviate:D
672 \__kernel_primitive:NN \pdfpageheight \tex_pageheight:D
673 \__kernel_primitive:NN \pdfpagewidth \tex_pagewidth:D
674 \__kernel_primitive:NN \pdfpkmode \tex_pkmode:D
675 \__kernel_primitive:NN \pdfpkresolution \tex_pkresolution:D
676 \__kernel_primitive:NN \pdfprimitive \tex_primitive:D
677 \__kernel_primitive:NN \pdfprotrudechars \tex_protrudechars:D
678 \__kernel_primitive:NN \pdfpxdimen \tex_pxdimen:D
679 \__kernel_primitive:NN \pdfrandomseed \tex_randomseed:D
680 \__kernel_primitive:NN \pdfresettimer \tex_resettimer:D
681 \__kernel_primitive:NN \pdfsavepos \tex_savepos:D
682 \__kernel_primitive:NN \pdfsetrandomseed \tex_setrandomseed:D
683 \__kernel_primitive:NN \pdfshellescape \tex_shellescape:D
684 \__kernel_primitive:NN \pdftracingfonts \tex_tracingfonts:D
685 \__kernel_primitive:NN \pdfuniformdeviate \tex_uniformdeviate:D

```

The version primitives are not related to PDF mode but are pdfTeX-specific, so again are carried forward unchanged.

```

686 \__kernel_primitive:NN \pdfTeXbanner \tex_pdfTeXbanner:D
687 \__kernel_primitive:NN \pdfTeXrevision \tex_pdfTeXrevision:D
688 \__kernel_primitive:NN \pdfTeXversion \tex_pdfTeXversion:D

```

These ones appear in pdfTeX but don't have pdf in the name at all: no decisions to make.

```

689 \__kernel_primitive:NN \efcode \tex_efcode:D
690 \__kernel_primitive:NN \ifincsname \tex_ifincsname:D
691 \__kernel_primitive:NN \leftmarginkern \tex_leftmarginkern:D
692 \__kernel_primitive:NN \letterspacefont \tex_letterspacefont:D
693 \__kernel_primitive:NN \lpcode \tex_lpcode:D
694 \__kernel_primitive:NN \quitvmode \tex_quitvmode:D
695 \__kernel_primitive:NN \rightmarginkern \tex_rightmarginkern:D
696 \__kernel_primitive:NN \rpcode \tex_rpcode:D
697 \__kernel_primitive:NN \synctex \tex_synctex:D
698 \__kernel_primitive:NN \tagcode \tex_tagcode:D

```

Post pdfTeX primitive availability gets more complex. Both XeTeX and LuaTeX have varying names for some primitives from pdfTeX. Particularly for LuaTeX tracking all of that would be hard. Instead, we now check that we only save primitives if they actually exist.

```

699 </names | package>
700 {*package>
701 \tex_long:D \tex_def:D \use_ii:nn #1#2 {#2}
702 \tex_long:D \tex_def:D \use_none:n #1 { }
703 \tex_long:D \tex_def:D \__kernel_primitive:NN #1#2
704 {

```

```

705     \tex_ifdefined:D #1
706     \tex_expandafter:D \use_ii:nn
707     \tex_fi:D
708     \use_none:n { \tex_global:D \tex_let:D #2 #1 }
709 }
710 \</package>
711 \<*names | package>

```

Some pdfTeX primitives are handled here because they got dropped in LuaTeX but the corresponding internal names are emulated later. The Lua code is already loaded at this point, so we shouldn't overwrite them.

```

712 \__kernel_primitive:NN \pdfstrcmp           \tex_strcmp:D
713 \__kernel_primitive:NN \pdffilesize         \tex_filesiz:D
714 \__kernel_primitive:NN \pdfmdfivesum       \tex_mdfivesum:D
715 \__kernel_primitive:NN \pdffilemoddate     \tex_filemoddate:D
716 \__kernel_primitive:NN \pdffiledump        \tex_filedump:D

```

X_YTeX-specific primitives. Note that X_YTeX's \strcmp is handled earlier and is “rolled up” into \pdfstrcmp. A few cross-compatibility names which lack the pdf of the original are handled later.

```

717 \__kernel_primitive:NN \suppressfontnotfounderror
718 \tex_suppressfontnotfounderror:D
719 \__kernel_primitive:NN \XeTeXcharclass      \tex_XeTeXcharclass:D
720 \__kernel_primitive:NN \XeTeXcharglyph     \tex_XeTeXcharglyph:D
721 \__kernel_primitive:NN \XeTeXcountfeatures \tex_XeTeXcountfeatures:D
722 \__kernel_primitive:NN \XeTeXcountglyphs   \tex_XeTeXcountglyphs:D
723 \__kernel_primitive:NN \XeTeXcountselectors \tex_XeTeXcountselectors:D
724 \__kernel_primitive:NN \XeTeXcountvariations \tex_XeTeXcountvariations:D
725 \__kernel_primitive:NN \XeTeXdefaultencoding \tex_XeTeXdefaultencoding:D
726 \__kernel_primitive:NN \XeTeXdashbreakstate \tex_XeTeXdashbreakstate:D
727 \__kernel_primitive:NN \XeTeXfeaturecode   \tex_XeTeXfeaturecode:D
728 \__kernel_primitive:NN \XeTeXfeaturename    \tex_XeTeXfeaturename:D
729 \__kernel_primitive:NN \XeTeXfindfeaturebyname
730 \tex_XeTeXfindfeaturebyname:D
731 \__kernel_primitive:NN \XeTeXfindselectorbyname
732 \tex_XeTeXfindselectorbyname:D
733 \__kernel_primitive:NN \XeTeXfindvariationbyname
734 \tex_XeTeXfindvariationbyname:D
735 \__kernel_primitive:NN \XeTeXfirstfontchar \tex_XeTeXfirstfontchar:D
736 \__kernel_primitive:NN \XeTeXfonttype      \tex_XeTeXfonttype:D
737 \__kernel_primitive:NN \XeTeXgenerateactualtext
738 \tex_XeTeXgenerateactualtext:D
739 \__kernel_primitive:NN \XeTeXglyph         \tex_XeTeXglyph:D
740 \__kernel_primitive:NN \XeTeXglyphbounds   \tex_XeTeXglyphbounds:D
741 \__kernel_primitive:NN \XeTeXglyphindex    \tex_XeTeXglyphindex:D
742 \__kernel_primitive:NN \XeTeXglyphname     \tex_XeTeXglyphname:D
743 \__kernel_primitive:NN \XeTeXinputencoding \tex_XeTeXinputencoding:D
744 \__kernel_primitive:NN \XeTeXinputnormalization
745 \tex_XeTeXinputnormalization:D
746 \__kernel_primitive:NN \XeTeXinterchartokenstate
747 \tex_XeTeXinterchartokenstate:D
748 \__kernel_primitive:NN \XeTeXinterchartoks \tex_XeTeXinterchartoks:D
749 \__kernel_primitive:NN \XeTeXisdefaultselector
750 \tex_XeTeXisdefaultselector:D
751 \__kernel_primitive:NN \XeTeXisexclusivefeature

```

752	<code>\tex_XeTeXisexclusivefeature:D</code>	
753	<code>__kernel_primitive:NN \XeTeXlastfontchar</code>	<code>\tex_XeTeXlastfontchar:D</code>
754	<code>__kernel_primitive:NN \XeTeXlinebreakskip</code>	<code>\tex_XeTeXlinebreakskip:D</code>
755	<code>__kernel_primitive:NN \XeTeXlinebreaklocale</code>	<code>\tex_XeTeXlinebreaklocale:D</code>
756	<code>__kernel_primitive:NN \XeTeXlinebreakpenalty</code>	<code>\tex_XeTeXlinebreakpenalty:D</code>
757	<code>__kernel_primitive:NN \XeTeXOTcountfeatures</code>	<code>\tex_XeTeXOTcountfeatures:D</code>
758	<code>__kernel_primitive:NN \XeTeXOTcountlanguages</code>	<code>\tex_XeTeXOTcountlanguages:D</code>
759	<code>__kernel_primitive:NN \XeTeXOTcountscripts</code>	<code>\tex_XeTeXOTcountscripts:D</code>
760	<code>__kernel_primitive:NN \XeTeXOTfeaturetag</code>	<code>\tex_XeTeXOTfeaturetag:D</code>
761	<code>__kernel_primitive:NN \XeTeXOTlanguage tag</code>	<code>\tex_XeTeXOTlanguage tag:D</code>
762	<code>__kernel_primitive:NN \XeTeXOTscripttag</code>	<code>\tex_XeTeXOTscripttag:D</code>
763	<code>__kernel_primitive:NN \XeTeXpdf file</code>	<code>\tex_XeTeXpdf file:D</code>
764	<code>__kernel_primitive:NN \XeTeXpdfpagecount</code>	<code>\tex_XeTeXpdfpagecount:D</code>
765	<code>__kernel_primitive:NN \XeTeXpic file</code>	<code>\tex_XeTeXpic file:D</code>
766	<code>__kernel_primitive:NN \XeTeXrevision</code>	<code>\tex_XeTeXrevision:D</code>
767	<code>__kernel_primitive:NN \XeTeXselectorname</code>	<code>\tex_XeTeXselectorname:D</code>
768	<code>__kernel_primitive:NN \XeTeXtracing fonts</code>	<code>\tex_XeTeXtracing fonts:D</code>
769	<code>__kernel_primitive:NN \XeTeXupwards mode</code>	<code>\tex_XeTeXupwards mode:D</code>
770	<code>__kernel_primitive:NN \XeTeXuse glyph metrics</code>	<code>\tex_XeTeXuse glyph metrics:D</code>
771	<code>__kernel_primitive:NN \XeTeXvariation</code>	<code>\tex_XeTeXvariation:D</code>
772	<code>__kernel_primitive:NN \XeTeXvariation on default</code>	<code>\tex_XeTeXvariation on default:D</code>
773	<code>__kernel_primitive:NN \XeTeXvariation max</code>	<code>\tex_XeTeXvariation max:D</code>
774	<code>__kernel_primitive:NN \XeTeXvariation min</code>	<code>\tex_XeTeXvariation min:D</code>
775	<code>__kernel_primitive:NN \XeTeXvariation name</code>	<code>\tex_XeTeXvariation name:D</code>
776	<code>__kernel_primitive:NN \XeTeXversion</code>	<code>\tex_XeTeXversion:D</code>

Primitives from pdfTeX that XeTeX renames: also helps with LuaTeX.

777	<code>__kernel_primitive:NN \creationdate</code>	<code>\tex_creationdate:D</code>
778	<code>__kernel_primitive:NN \elapsedtime</code>	<code>\tex_elapsedtime:D</code>
779	<code>__kernel_primitive:NN \filedump</code>	<code>\tex_filedump:D</code>
780	<code>__kernel_primitive:NN \filemoddate</code>	<code>\tex_filemoddate:D</code>
781	<code>__kernel_primitive:NN \filesize</code>	<code>\tex_filesize:D</code>
782	<code>__kernel_primitive:NN \mdfivesum</code>	<code>\tex_mdfivesum:D</code>
783	<code>__kernel_primitive:NN \ifprimitive</code>	<code>\tex_ifprimitive:D</code>
784	<code>__kernel_primitive:NN \primitive</code>	<code>\tex_primitive:D</code>
785	<code>__kernel_primitive:NN \resettimer</code>	<code>\tex_resettimer:D</code>
786	<code>__kernel_primitive:NN \shellescape</code>	<code>\tex_shellescape:D</code>

Primitives from LuaTeX, some of which have been ported back to XeTeX.

787	<code>__kernel_primitive:NN \alignmark</code>	<code>\tex_alignmark:D</code>
788	<code>__kernel_primitive:NN \align tab</code>	<code>\tex_align tab:D</code>
789	<code>__kernel_primitive:NN \attribute</code>	<code>\tex_attribute:D</code>
790	<code>__kernel_primitive:NN \attributedef</code>	<code>\tex_attributedef:D</code>
791	<code>__kernel_primitive:NN \automaticdiscretionary</code>	
792	<code>\tex_automaticdiscretionary:D</code>	
793	<code>__kernel_primitive:NN \automatichyphenmode</code>	<code>\tex_automatichyphenmode:D</code>
794	<code>__kernel_primitive:NN \automatichyphenpenalty</code>	
795	<code>\tex_automatichyphenpenalty:D</code>	
796	<code>__kernel_primitive:NN \beginsname</code>	<code>\tex_beginsname:D</code>
797	<code>__kernel_primitive:NN \bodydir</code>	<code>\tex_bodydir:D</code>
798	<code>__kernel_primitive:NN \bodydirection</code>	<code>\tex_bodydirection:D</code>
799	<code>__kernel_primitive:NN \boxdir</code>	<code>\tex_boxdir:D</code>
800	<code>__kernel_primitive:NN \boxdirection</code>	<code>\tex_boxdirection:D</code>
801	<code>__kernel_primitive:NN \breakafterdirmode</code>	<code>\tex_breakafterdirmode:D</code>
802	<code>__kernel_primitive:NN \catcodetable</code>	<code>\tex_catcodetable:D</code>

803	<code>__kernel_primitive:NN \clearmarks</code>	<code>\tex_clearmarks:D</code>
804	<code>__kernel_primitive:NN \crampeddisplaystyle</code>	<code>\tex_crampeddisplaystyle:D</code>
805	<code>__kernel_primitive:NN \crampedscriptscriptstyle</code>	
806	<code>\tex_crampedscriptscriptstyle:D</code>	
807	<code>__kernel_primitive:NN \crampedscriptstyle</code>	<code>\tex_crampedscriptstyle:D</code>
808	<code>__kernel_primitive:NN \crampedtextstyle</code>	<code>\tex_crampedtextstyle:D</code>
809	<code>__kernel_primitive:NN \csstring</code>	<code>\tex_csstring:D</code>
810	<code>__kernel_primitive:NN \directlua</code>	<code>\tex_directlua:D</code>
811	<code>__kernel_primitive:NN \dviextension</code>	<code>\tex_dviextension:D</code>
812	<code>__kernel_primitive:NN \dvifedback</code>	<code>\tex_dvifedback:D</code>
813	<code>__kernel_primitive:NN \dvivariable</code>	<code>\tex_dvivariable:D</code>
814	<code>__kernel_primitive:NN \eTeXglueshrinkorder</code>	<code>\tex_eTeXglueshrinkorder:D</code>
815	<code>__kernel_primitive:NN \eTeXgluestretchorder</code>	<code>\tex_eTeXgluestretchorder:D</code>
816	<code>__kernel_primitive:NN \etoksapp</code>	<code>\tex_etoksapp:D</code>
817	<code>__kernel_primitive:NN \etokspre</code>	<code>\tex_etokspre:D</code>
818	<code>__kernel_primitive:NN \exceptionpenalty</code>	<code>\tex_exceptionpenalty:D</code>
819	<code>__kernel_primitive:NN \explicithyphenpenalty</code>	<code>\tex_explicithyphenpenalty:D</code>
820	<code>__kernel_primitive:NN \expanded</code>	<code>\tex_expanded:D</code>
821	<code>__kernel_primitive:NN \explicitdiscretionary</code>	<code>\tex_explicitdiscretionary:D</code>
822	<code>__kernel_primitive:NN \firstvalidlanguage</code>	<code>\tex_firstvalidlanguage:D</code>
823	<code>__kernel_primitive:NN \fontid</code>	<code>\tex_fontid:D</code>
824	<code>__kernel_primitive:NN \formatname</code>	<code>\tex_formatname:D</code>
825	<code>__kernel_primitive:NN \hrcode</code>	<code>\tex_hrcode:D</code>
826	<code>__kernel_primitive:NN \hpack</code>	<code>\tex_hpack:D</code>
827	<code>__kernel_primitive:NN \hyphenationbounds</code>	<code>\tex_hyphenationbounds:D</code>
828	<code>__kernel_primitive:NN \hyphenationmin</code>	<code>\tex_hyphenationmin:D</code>
829	<code>__kernel_primitive:NN \hyphenpenaltymode</code>	<code>\tex_hyphenpenaltymode:D</code>
830	<code>__kernel_primitive:NN \gleaders</code>	<code>\tex_gleaders:D</code>
831	<code>__kernel_primitive:NN \ifcondition</code>	<code>\tex_ifcondition:D</code>
832	<code>__kernel_primitive:NN \immediateassigned</code>	<code>\tex_immediateassigned:D</code>
833	<code>__kernel_primitive:NN \immediateassignment</code>	<code>\tex_immediateassignment:D</code>
834	<code>__kernel_primitive:NN \initcatcodetable</code>	<code>\tex_initcatcodetable:D</code>
835	<code>__kernel_primitive:NN \lastnamedcs</code>	<code>\tex_lastnamedcs:D</code>
836	<code>__kernel_primitive:NN \latelua</code>	<code>\tex_latelua:D</code>
837	<code>__kernel_primitive:NN \lateluafunction</code>	<code>\tex_lateluafunction:D</code>
838	<code>__kernel_primitive:NN \leftghost</code>	<code>\tex_leftghost:D</code>
839	<code>__kernel_primitive:NN \letcharcode</code>	<code>\tex_letcharcode:D</code>
840	<code>__kernel_primitive:NN \linedir</code>	<code>\tex_linedir:D</code>
841	<code>__kernel_primitive:NN \linedirection</code>	<code>\tex_linedirection:D</code>
842	<code>__kernel_primitive:NN \localbrokenpenalty</code>	<code>\tex_localbrokenpenalty:D</code>
843	<code>__kernel_primitive:NN \localinterlinepenalty</code>	<code>\tex_localinterlinepenalty:D</code>
844	<code>__kernel_primitive:NN \luabytecode</code>	<code>\tex_luabytecode:D</code>
845	<code>__kernel_primitive:NN \luabytecodecall</code>	<code>\tex_luabytecodecall:D</code>
846	<code>__kernel_primitive:NN \luacopyinputnodes</code>	<code>\tex_luacopyinputnodes:D</code>
847	<code>__kernel_primitive:NN \luaodef</code>	<code>\tex_luaodef:D</code>
848	<code>__kernel_primitive:NN \lcalleftbox</code>	<code>\tex_lcalleftbox:D</code>
849	<code>__kernel_primitive:NN \lcalrightbox</code>	<code>\tex_lcalrightbox:D</code>
850	<code>__kernel_primitive:NN \luaescapestring</code>	<code>\tex_luaescapestring:D</code>
851	<code>__kernel_primitive:NN \luafunction</code>	<code>\tex_luafunction:D</code>
852	<code>__kernel_primitive:NN \luafunctioncall</code>	<code>\tex_luafunctioncall:D</code>
853	<code>__kernel_primitive:NN \luatexbanner</code>	<code>\tex_luatexbanner:D</code>
854	<code>__kernel_primitive:NN \luatexrevision</code>	<code>\tex_luatexrevision:D</code>
855	<code>__kernel_primitive:NN \luatexversion</code>	<code>\tex_luatexversion:D</code>
856	<code>__kernel_primitive:NN \mathdelimitersmode</code>	<code>\tex_mathdelimitersmode:D</code>

857	<code>__kernel_primitive:NN \mathdir</code>	<code>\tex_mathdir:D</code>
858	<code>__kernel_primitive:NN \mathdirection</code>	<code>\tex_mathdirection:D</code>
859	<code>__kernel_primitive:NN \mathdisplayskipmode</code>	<code>\tex_mathdisplayskipmode:D</code>
860	<code>__kernel_primitive:NN \matheqnogapstep</code>	<code>\tex_matheqnogapstep:D</code>
861	<code>__kernel_primitive:NN \mathnolimitsmode</code>	<code>\tex_mathnolimitsmode:D</code>
862	<code>__kernel_primitive:NN \mathoption</code>	<code>\tex_mathoption:D</code>
863	<code>__kernel_primitive:NN \mathpenaltiesmode</code>	<code>\tex_mathpenaltiesmode:D</code>
864	<code>__kernel_primitive:NN \mathrulesfam</code>	<code>\tex_mathrulesfam:D</code>
865	<code>__kernel_primitive:NN \mathscriptsmode</code>	<code>\tex_mathscriptsmode:D</code>
866	<code>__kernel_primitive:NN \mathscriptboxmode</code>	<code>\tex_mathscriptboxmode:D</code>
867	<code>__kernel_primitive:NN \mathscriptcharmode</code>	<code>\tex_mathscriptcharmode:D</code>
868	<code>__kernel_primitive:NN \mathstyle</code>	<code>\tex_mathstyle:D</code>
869	<code>__kernel_primitive:NN \mathsurroundmode</code>	<code>\tex_mathsurroundmode:D</code>
870	<code>__kernel_primitive:NN \mathsurroundskip</code>	<code>\tex_mathsurroundskip:D</code>
871	<code>__kernel_primitive:NN \nohrule</code>	<code>\tex_nohrule:D</code>
872	<code>__kernel_primitive:NN \nokerns</code>	<code>\tex_nokerns:D</code>
873	<code>__kernel_primitive:NN \noligs</code>	<code>\tex_noligs:D</code>
874	<code>__kernel_primitive:NN \nospaces</code>	<code>\tex_nospaces:D</code>
875	<code>__kernel_primitive:NN \novrule</code>	<code>\tex_novrule:D</code>
876	<code>__kernel_primitive:NN \outputbox</code>	<code>\tex_outputbox:D</code>
877	<code>__kernel_primitive:NN \pagebottomoffset</code>	<code>\tex_pagebottomoffset:D</code>
878	<code>__kernel_primitive:NN \pagedir</code>	<code>\tex_pagedir:D</code>
879	<code>__kernel_primitive:NN \pagedirection</code>	<code>\tex_pagedirection:D</code>
880	<code>__kernel_primitive:NN \pageleftoffset</code>	<code>\tex_pageleftoffset:D</code>
881	<code>__kernel_primitive:NN \pagerightoffset</code>	<code>\tex_pagerightoffset:D</code>
882	<code>__kernel_primitive:NN \pagetopoffset</code>	<code>\tex_pagetopoffset:D</code>
883	<code>__kernel_primitive:NN \paddir</code>	<code>\tex_paddir:D</code>
884	<code>__kernel_primitive:NN \pardirection</code>	<code>\tex_pardirection:D</code>
885	<code>__kernel_primitive:NN \pdfextension</code>	<code>\tex_pdfextension:D</code>
886	<code>__kernel_primitive:NN \pdffeedback</code>	<code>\tex_pdffeedback:D</code>
887	<code>__kernel_primitive:NN \pdfvariable</code>	<code>\tex_pdfvariable:D</code>
888	<code>__kernel_primitive:NN \postexhyphenchar</code>	<code>\tex_postexhyphenchar:D</code>
889	<code>__kernel_primitive:NN \posthyphenchar</code>	<code>\tex_posthyphenchar:D</code>
890	<code>__kernel_primitive:NN \prebinoppenalty</code>	<code>\tex_prebinoppenalty:D</code>
891	<code>__kernel_primitive:NN \predisplaygapfactor</code>	<code>\tex_predisplaygapfactor:D</code>
892	<code>__kernel_primitive:NN \preexhyphenchar</code>	<code>\tex_preexhyphenchar:D</code>
893	<code>__kernel_primitive:NN \prehyphenchar</code>	<code>\tex_prehyphenchar:D</code>
894	<code>__kernel_primitive:NN \prerelpenalty</code>	<code>\tex_prerelpenalty:D</code>
895	<code>__kernel_primitive:NN \rightghost</code>	<code>\tex_rightghost:D</code>
896	<code>__kernel_primitive:NN \savecatcodetable</code>	<code>\tex_savecatcodetable:D</code>
897	<code>__kernel_primitive:NN \scantextokens</code>	<code>\tex_scantextokens:D</code>
898	<code>__kernel_primitive:NN \setfontid</code>	<code>\tex_setfontid:D</code>
899	<code>__kernel_primitive:NN \shapemode</code>	<code>\tex_shapemode:D</code>
900	<code>__kernel_primitive:NN \suppressifcsnameerror</code>	<code>\tex_suppressifcsnameerror:D</code>
901	<code>__kernel_primitive:NN \suppresslongerror</code>	<code>\tex_suppresslongerror:D</code>
902	<code>__kernel_primitive:NN \suppressmathparerror</code>	<code>\tex_suppressmathparerror:D</code>
903	<code>__kernel_primitive:NN \suppressoutererror</code>	<code>\tex_suppressoutererror:D</code>
904	<code>__kernel_primitive:NN \suppressprimitiveerror</code>	
905	<code>\tex_suppressprimitiveerror:D</code>	
906	<code>__kernel_primitive:NN \textdir</code>	<code>\tex_textdir:D</code>
907	<code>__kernel_primitive:NN \textdirection</code>	<code>\tex_textdirection:D</code>
908	<code>__kernel_primitive:NN \toksapp</code>	<code>\tex_toksapp:D</code>
909	<code>__kernel_primitive:NN \tokspre</code>	<code>\tex_tokspre:D</code>
910	<code>__kernel_primitive:NN \tpack</code>	<code>\tex_tpack:D</code>

911 `__kernel_primitive:NN \vpack` `\tex_vpack:D`

Primitives from pdfTeX that LuaTeX renames.

912 `__kernel_primitive:NN \adjustspacing` `\tex_adjustspacing:D`
913 `__kernel_primitive:NN \copyfont` `\tex_copyfont:D`
914 `__kernel_primitive:NN \draftmode` `\tex_draftmode:D`
915 `__kernel_primitive:NN \expandglyphsinfont` `\tex_fontexpand:D`
916 `__kernel_primitive:NN \ifabsdim` `\tex_ifabsdim:D`
917 `__kernel_primitive:NN \ifabsnum` `\tex_ifabsnum:D`
918 `__kernel_primitive:NN \ignoreligaturesinfont` `\tex_ignoreligaturesinfont:D`
919 `__kernel_primitive:NN \insertht` `\tex_insertht:D`
920 `__kernel_primitive:NN \lastsavedboxresourceindex`
921 `\tex_pdflastxform:D`
922 `__kernel_primitive:NN \lastsavedimageresourceindex`
923 `\tex_pdflastximage:D`
924 `__kernel_primitive:NN \lastsavedimageresourcepages`
925 `\tex_pdflastximagepages:D`
926 `__kernel_primitive:NN \lastxpos` `\tex_lastxpos:D`
927 `__kernel_primitive:NN \lastypos` `\tex_lastypos:D`
928 `__kernel_primitive:NN \normaldeviate` `\tex_normaldeviate:D`
929 `__kernel_primitive:NN \outputmode` `\tex_pdfoutput:D`
930 `__kernel_primitive:NN \pageheight` `\tex_pageheight:D`
931 `__kernel_primitive:NN \pagewidth` `\tex_pagewidth:D`
932 `__kernel_primitive:NN \protrudechars` `\tex_protrudechars:D`
933 `__kernel_primitive:NN \pxdimen` `\tex_pxdimen:D`
934 `__kernel_primitive:NN \randomseed` `\tex_randomseed:D`
935 `__kernel_primitive:NN \useboxresource` `\tex_pdfrefxform:D`
936 `__kernel_primitive:NN \useimageresource` `\tex_pdfrefximage:D`
937 `__kernel_primitive:NN \savepos` `\tex_savepos:D`
938 `__kernel_primitive:NN \saveboxresource` `\tex_pdfxform:D`
939 `__kernel_primitive:NN \saveimageresource` `\tex_pdfximage:D`
940 `__kernel_primitive:NN \setrandomseed` `\tex_setrandomseed:D`
941 `__kernel_primitive:NN \tracingfonts` `\tex_tracingfonts:D`
942 `__kernel_primitive:NN \uniformdeviate` `\tex_uniformdeviate:D`

The set of Unicode math primitives were introduced by XeTeX and LuaTeX in a somewhat complex fashion: a few first as `\XeTeX...` which were then renamed with LuaTeX having a lot more. These names now all start `\U...` and mainly `\Umath...`.

943 `__kernel_primitive:NN \Uchar` `\tex_Uchar:D`
944 `__kernel_primitive:NN \Ucharcat` `\tex_Ucharcat:D`
945 `__kernel_primitive:NN \Udelcode` `\tex_Udelcode:D`
946 `__kernel_primitive:NN \Udelcodenum` `\tex_Udelcodenum:D`
947 `__kernel_primitive:NN \Udelimiter` `\tex_Udelimiter:D`
948 `__kernel_primitive:NN \Udelimiterover` `\tex_Udelimiterover:D`
949 `__kernel_primitive:NN \Udelimiterunder` `\tex_Udelimiterunder:D`
950 `__kernel_primitive:NN \Uxextensible` `\tex_Uxextensible:D`
951 `__kernel_primitive:NN \Umathaccent` `\tex_Umathaccent:D`
952 `__kernel_primitive:NN \Umathaxis` `\tex_Umathaxis:D`
953 `__kernel_primitive:NN \Umathbinbinspacing` `\tex_Umathbinbinspacing:D`
954 `__kernel_primitive:NN \Umathbinclonespacing` `\tex_Umathbinclonespacing:D`
955 `__kernel_primitive:NN \Umathbininnerspacing` `\tex_Umathbininnerspacing:D`
956 `__kernel_primitive:NN \Umathbinopenspacing` `\tex_Umathbinopenspacing:D`
957 `__kernel_primitive:NN \Umathbinopspacing` `\tex_Umathbinopspacing:D`
958 `__kernel_primitive:NN \Umathbinordspacing` `\tex_Umathbinordspacing:D`
959 `__kernel_primitive:NN \Umathbinpunctspacing` `\tex_Umathbinpunctspacing:D`

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960 \__kernel_primitive:NN \Umathbinrelspacing \tex_Umathbinrelspacing:D
961 \__kernel_primitive:NN \Umathchar \tex_Umathchar:D
962 \__kernel_primitive:NN \Umathcharclass \tex_Umathcharclass:D
963 \__kernel_primitive:NN \Umathchardef \tex_Umathchardef:D
964 \__kernel_primitive:NN \Umathcharfam \tex_Umathcharfam:D
965 \__kernel_primitive:NN \Umathcharnum \tex_Umathcharnum:D
966 \__kernel_primitive:NN \Umathcharnumdef \tex_Umathcharnumdef:D
967 \__kernel_primitive:NN \Umathcharslot \tex_Umathcharslot:D
968 \__kernel_primitive:NN \Umathclosebinspacing \tex_Umathclosebinspacing:D
969 \__kernel_primitive:NN \Umathcloseclosespacing
970 \tex_Umathcloseclosespacing:D
971 \__kernel_primitive:NN \Umathcloseinnerspacing
972 \tex_Umathcloseinnerspacing:D
973 \__kernel_primitive:NN \Umathcloseopenspacing \tex_Umathcloseopenspacing:D
974 \__kernel_primitive:NN \Umathcloseopspacing \tex_Umathcloseopspacing:D
975 \__kernel_primitive:NN \Umathcloseordspacing \tex_Umathcloseordspacing:D
976 \__kernel_primitive:NN \Umathclosepunctspacing
977 \tex_Umathclosepunctspacing:D
978 \__kernel_primitive:NN \Umathcloserelspacing \tex_Umathcloserelspacing:D
979 \__kernel_primitive:NN \Umathcode \tex_Umathcode:D
980 \__kernel_primitive:NN \Umathcodenum \tex_Umathcodenum:D
981 \__kernel_primitive:NN \Umathconnectoroverlapmin
982 \tex_Umathconnectoroverlapmin:D
983 \__kernel_primitive:NN \Umathfractiondelsize \tex_Umathfractiondelsize:D
984 \__kernel_primitive:NN \Umathfractiondenomdown
985 \tex_Umathfractiondenomdown:D
986 \__kernel_primitive:NN \Umathfractiondenomvgap
987 \tex_Umathfractiondenomvgap:D
988 \__kernel_primitive:NN \Umathfractionnumup \tex_Umathfractionnumup:D
989 \__kernel_primitive:NN \Umathfractionnumvgap \tex_Umathfractionnumvgap:D
990 \__kernel_primitive:NN \Umathfractionrule \tex_Umathfractionrule:D
991 \__kernel_primitive:NN \Umathinnerbinspacing \tex_Umathinnerbinspacing:D
992 \__kernel_primitive:NN \Umathinnerclosespacing
993 \tex_Umathinnerclosespacing:D
994 \__kernel_primitive:NN \Umathinnerinnerspacing
995 \tex_Umathinnerinnerspacing:D
996 \__kernel_primitive:NN \Umathinneropenspacing \tex_Umathinneropenspacing:D
997 \__kernel_primitive:NN \Umathinneropspacing \tex_Umathinneropspacing:D
998 \__kernel_primitive:NN \Umathinnerordspacing \tex_Umathinnerordspacing:D
999 \__kernel_primitive:NN \Umathinnerpunctspacing
1000 \tex_Umathinnerpunctspacing:D
1001 \__kernel_primitive:NN \Umathinnerrelspacing \tex_Umathinnerrelspacing:D
1002 \__kernel_primitive:NN \Umathlimitabovebgap \tex_Umathlimitabovebgap:D
1003 \__kernel_primitive:NN \Umathlimitabovekern \tex_Umathlimitabovekern:D
1004 \__kernel_primitive:NN \Umathlimitabovevgap \tex_Umathlimitabovevgap:D
1005 \__kernel_primitive:NN \Umathlimitbelowbgap \tex_Umathlimitbelowbgap:D
1006 \__kernel_primitive:NN \Umathlimitbelowkern \tex_Umathlimitbelowkern:D
1007 \__kernel_primitive:NN \Umathlimitbelowvgap \tex_Umathlimitbelowvgap:D
1008 \__kernel_primitive:NN \Umathnolimitsubfactor \tex_Umathnolimitsubfactor:D
1009 \__kernel_primitive:NN \Umathnolimitsupfactor \tex_Umathnolimitsupfactor:D
1010 \__kernel_primitive:NN \Umathopbinspacing \tex_Umathopbinspacing:D
1011 \__kernel_primitive:NN \Umathopclosespacing \tex_Umathopclosespacing:D
1012 \__kernel_primitive:NN \Umathopenbinspacing \tex_Umathopenbinspacing:D
1013 \__kernel_primitive:NN \Umathopenclosespacing \tex_Umathopenclosespacing:D

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1014 \__kernel_primitive:NN \Umathopeninnerspacing \tex_Umathopeninnerspacing:D
1015 \__kernel_primitive:NN \Umathopenopenspacing \tex_Umathopenopenspacing:D
1016 \__kernel_primitive:NN \Umathopenopspacing \tex_Umathopenopspacing:D
1017 \__kernel_primitive:NN \Umathopenordspacing \tex_Umathopenordspacing:D
1018 \__kernel_primitive:NN \Umathopenpunctspacing \tex_Umathopenpunctspacing:D
1019 \__kernel_primitive:NN \Umathopenrelspacing \tex_Umathopenrelspacing:D
1020 \__kernel_primitive:NN \Umathoperatorsize \tex_Umathoperatorsize:D
1021 \__kernel_primitive:NN \Umathopinnerspacing \tex_Umathopinnerspacing:D
1022 \__kernel_primitive:NN \Umathopopenspacing \tex_Umathopopenspacing:D
1023 \__kernel_primitive:NN \Umathopopspacing \tex_Umathopopspacing:D
1024 \__kernel_primitive:NN \Umathopordspacing \tex_Umathopordspacing:D
1025 \__kernel_primitive:NN \Umathoppunctspacing \tex_Umathoppunctspacing:D
1026 \__kernel_primitive:NN \Umathoprelspacing \tex_Umathoprelspacing:D
1027 \__kernel_primitive:NN \Umathordbinspacing \tex_Umathordbinspacing:D
1028 \__kernel_primitive:NN \Umathordclosespacing \tex_Umathordclosespacing:D
1029 \__kernel_primitive:NN \Umathordinnerspacing \tex_Umathordinnerspacing:D
1030 \__kernel_primitive:NN \Umathordopenspacing \tex_Umathordopenspacing:D
1031 \__kernel_primitive:NN \Umathordopspacing \tex_Umathordopspacing:D
1032 \__kernel_primitive:NN \Umathordordspacing \tex_Umathordordspacing:D
1033 \__kernel_primitive:NN \Umathordpunctspacing \tex_Umathordpunctspacing:D
1034 \__kernel_primitive:NN \Umathordrelspacing \tex_Umathordrelspacing:D
1035 \__kernel_primitive:NN \Umathoverbarkern \tex_Umathoverbarkern:D
1036 \__kernel_primitive:NN \Umathoverbarrule \tex_Umathoverbarrule:D
1037 \__kernel_primitive:NN \Umathoverbarvgap \tex_Umathoverbarvgap:D
1038 \__kernel_primitive:NN \Umathoverdelimiterbgap
1039 \tex_Umathoverdelimiterbgap:D
1040 \__kernel_primitive:NN \Umathoverdelimitervgap
1041 \tex_Umathoverdelimitervgap:D
1042 \__kernel_primitive:NN \Umathpunctbinspacing \tex_Umathpunctbinspacing:D
1043 \__kernel_primitive:NN \Umathpunctclosespacing
1044 \tex_Umathpunctclosespacing:D
1045 \__kernel_primitive:NN \Umathpunctinnerspacing
1046 \tex_Umathpunctinnerspacing:D
1047 \__kernel_primitive:NN \Umathpunctopenspacing \tex_Umathpunctopenspacing:D
1048 \__kernel_primitive:NN \Umathpunctopspacing \tex_Umathpunctopspacing:D
1049 \__kernel_primitive:NN \Umathpunctordspacing \tex_Umathpunctordspacing:D
1050 \__kernel_primitive:NN \Umathpunctpunctspacing
1051 \tex_Umathpunctpunctspacing:D
1052 \__kernel_primitive:NN \Umathpunctrelspacing \tex_Umathpunctrelspacing:D
1053 \__kernel_primitive:NN \Umathquad \tex_Umathquad:D
1054 \__kernel_primitive:NN \Umathradicaldegreeafter
1055 \tex_Umathradicaldegreeafter:D
1056 \__kernel_primitive:NN \Umathradicaldegreebefore
1057 \tex_Umathradicaldegreebefore:D
1058 \__kernel_primitive:NN \Umathradicaldegreeraise
1059 \tex_Umathradicaldegreeraise:D
1060 \__kernel_primitive:NN \Umathradicalkern \tex_Umathradicalkern:D
1061 \__kernel_primitive:NN \Umathradicalrule \tex_Umathradicalrule:D
1062 \__kernel_primitive:NN \Umathradicalvgap \tex_Umathradicalvgap:D
1063 \__kernel_primitive:NN \Umathrelbinspacing \tex_Umathrelbinspacing:D
1064 \__kernel_primitive:NN \Umathrelclosespacing \tex_Umathrelclosespacing:D
1065 \__kernel_primitive:NN \Umathrelinnerspacing \tex_Umathrelinnerspacing:D
1066 \__kernel_primitive:NN \Umathrelopenspacing \tex_Umathrelopenspacing:D
1067 \__kernel_primitive:NN \Umathrelopspacing \tex_Umathrelopspacing:D

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1068	<code>_kernel_primitive:NN \Umathrelordspacing</code>	<code>\tex_Umathrelordspacing:D</code>
1069	<code>_kernel_primitive:NN \Umathrelpunctspacing</code>	<code>\tex_Umathrelpunctspacing:D</code>
1070	<code>_kernel_primitive:NN \Umathrelrelspacing</code>	<code>\tex_Umathrelrelspacing:D</code>
1071	<code>_kernel_primitive:NN \Umathskewedfractionhgap</code>	
1072	<code>\tex_Umathskewedfractionhgap:D</code>	
1073	<code>_kernel_primitive:NN \Umathskewedfractionvgap</code>	
1074	<code>\tex_Umathskewedfractionvgap:D</code>	
1075	<code>_kernel_primitive:NN \Umathspaceafterscript</code>	<code>\tex_Umathspaceafterscript:D</code>
1076	<code>_kernel_primitive:NN \Umathstackdenomdown</code>	<code>\tex_Umathstackdenomdown:D</code>
1077	<code>_kernel_primitive:NN \Umathstacknumup</code>	<code>\tex_Umathstacknumup:D</code>
1078	<code>_kernel_primitive:NN \Umathstackvgap</code>	<code>\tex_Umathstackvgap:D</code>
1079	<code>_kernel_primitive:NN \Umathsubshiftdown</code>	<code>\tex_Umathsubshiftdown:D</code>
1080	<code>_kernel_primitive:NN \Umathsubshiftdrop</code>	<code>\tex_Umathsubshiftdrop:D</code>
1081	<code>_kernel_primitive:NN \Umathsubsupshiftdown</code>	<code>\tex_Umathsubsupshiftdown:D</code>
1082	<code>_kernel_primitive:NN \Umathsubsupvgap</code>	<code>\tex_Umathsubsupvgap:D</code>
1083	<code>_kernel_primitive:NN \Umathsubtopmax</code>	<code>\tex_Umathsubtopmax:D</code>
1084	<code>_kernel_primitive:NN \Umathsupbottommin</code>	<code>\tex_Umathsupbottommin:D</code>
1085	<code>_kernel_primitive:NN \Umathsupshiftdrop</code>	<code>\tex_Umathsupshiftdrop:D</code>
1086	<code>_kernel_primitive:NN \Umathsupshiftup</code>	<code>\tex_Umathsupshiftup:D</code>
1087	<code>_kernel_primitive:NN \Umathsupsubbottommax</code>	<code>\tex_Umathsupsubbottommax:D</code>
1088	<code>_kernel_primitive:NN \Umathunderbarkern</code>	<code>\tex_Umathunderbarkern:D</code>
1089	<code>_kernel_primitive:NN \Umathunderbarrule</code>	<code>\tex_Umathunderbarrule:D</code>
1090	<code>_kernel_primitive:NN \Umathunderbarvgap</code>	<code>\tex_Umathunderbarvgap:D</code>
1091	<code>_kernel_primitive:NN \Umathunderdelimiterbgap</code>	
1092	<code>\tex_Umathunderdelimiterbgap:D</code>	
1093	<code>_kernel_primitive:NN \Umathunderdelimitervgap</code>	
1094	<code>\tex_Umathunderdelimitervgap:D</code>	
1095	<code>_kernel_primitive:NN \Unosubscript</code>	<code>\tex_Unosubscript:D</code>
1096	<code>_kernel_primitive:NN \Unosuperscript</code>	<code>\tex_Unosuperscript:D</code>
1097	<code>_kernel_primitive:NN \Uoverdelimiter</code>	<code>\tex_Uoverdelimiter:D</code>
1098	<code>_kernel_primitive:NN \Uradical</code>	<code>\tex_Uradical:D</code>
1099	<code>_kernel_primitive:NN \Uroot</code>	<code>\tex_Uroot:D</code>
1100	<code>_kernel_primitive:NN \Uskewed</code>	<code>\tex_Uskewed:D</code>
1101	<code>_kernel_primitive:NN \Uskewedwithdelims</code>	<code>\tex_Uskewedwithdelims:D</code>
1102	<code>_kernel_primitive:NN \Ustack</code>	<code>\tex_Ustack:D</code>
1103	<code>_kernel_primitive:NN \Ustartdisplaymath</code>	<code>\tex_Ustartdisplaymath:D</code>
1104	<code>_kernel_primitive:NN \Ustartmath</code>	<code>\tex_Ustartmath:D</code>
1105	<code>_kernel_primitive:NN \Ustopdisplaymath</code>	<code>\tex_Ustopdisplaymath:D</code>
1106	<code>_kernel_primitive:NN \Ustopmath</code>	<code>\tex_Ustopmath:D</code>
1107	<code>_kernel_primitive:NN \Usubscript</code>	<code>\tex_Usubscript:D</code>
1108	<code>_kernel_primitive:NN \Usuperscript</code>	<code>\tex_Usuperscript:D</code>
1109	<code>_kernel_primitive:NN \Uunderdelimiter</code>	<code>\tex_Uunderdelimiter:D</code>
1110	<code>_kernel_primitive:NN \Uvextensible</code>	<code>\tex_Uvextensible:D</code>

Primitives from pTeX.

1111	<code>_kernel_primitive:NN \autospaceing</code>	<code>\tex_autospaceing:D</code>
1112	<code>_kernel_primitive:NN \autoxspaceing</code>	<code>\tex_autoxspaceing:D</code>
1113	<code>_kernel_primitive:NN \currentcjktoken</code>	<code>\tex_currentcjktoken:D</code>
1114	<code>_kernel_primitive:NN \currentspacingmode</code>	<code>\tex_currentspacingmode:D</code>
1115	<code>_kernel_primitive:NN \currentxspacingmode</code>	<code>\tex_currentxspacingmode:D</code>
1116	<code>_kernel_primitive:NN \disinhibitglue</code>	<code>\tex_disinhibitglue:D</code>
1117	<code>_kernel_primitive:NN \dtou</code>	<code>\tex_dtou:D</code>
1118	<code>_kernel_primitive:NN \epTeXinputencoding</code>	<code>\tex_epTeXinputencoding:D</code>
1119	<code>_kernel_primitive:NN \epTeXversion</code>	<code>\tex_epTeXversion:D</code>
1120	<code>_kernel_primitive:NN \euc</code>	<code>\tex_euc:D</code>

1121	<code>__kernel_primitive:NN \hfi</code>	<code>\tex_hfi:D</code>
1122	<code>__kernel_primitive:NN \ifdbbox</code>	<code>\tex_ifdbbox:D</code>
1123	<code>__kernel_primitive:NN \ifddir</code>	<code>\tex_ifddir:D</code>
1124	<code>__kernel_primitive:NN \ifjfont</code>	<code>\tex_ifjfont:D</code>
1125	<code>__kernel_primitive:NN \ifmbox</code>	<code>\tex_ifmbox:D</code>
1126	<code>__kernel_primitive:NN \ifmdir</code>	<code>\tex_ifmdir:D</code>
1127	<code>__kernel_primitive:NN \iftbox</code>	<code>\tex_iftbox:D</code>
1128	<code>__kernel_primitive:NN \iftfont</code>	<code>\tex_iftfont:D</code>
1129	<code>__kernel_primitive:NN \iftdir</code>	<code>\tex_iftdir:D</code>
1130	<code>__kernel_primitive:NN \ifybox</code>	<code>\tex_ifybox:D</code>
1131	<code>__kernel_primitive:NN \ifydir</code>	<code>\tex_ifydir:D</code>
1132	<code>__kernel_primitive:NN \inhibitglue</code>	<code>\tex_inhibitglue:D</code>
1133	<code>__kernel_primitive:NN \inhibitxspcode</code>	<code>\tex_inhibitxspcode:D</code>
1134	<code>__kernel_primitive:NN \jcharwidowpenalty</code>	<code>\tex_jcharwidowpenalty:D</code>
1135	<code>__kernel_primitive:NN \jfam</code>	<code>\tex_jfam:D</code>
1136	<code>__kernel_primitive:NN \jfont</code>	<code>\tex_jfont:D</code>
1137	<code>__kernel_primitive:NN \jis</code>	<code>\tex_jis:D</code>
1138	<code>__kernel_primitive:NN \kanjiskip</code>	<code>\tex_kanjiskip:D</code>
1139	<code>__kernel_primitive:NN \kansuji</code>	<code>\tex_kansuji:D</code>
1140	<code>__kernel_primitive:NN \kansujichar</code>	<code>\tex_kansujichar:D</code>
1141	<code>__kernel_primitive:NN \kcatcode</code>	<code>\tex_kcatcode:D</code>
1142	<code>__kernel_primitive:NN \kuten</code>	<code>\tex_kuten:D</code>
1143	<code>__kernel_primitive:NN \lastnodechar</code>	<code>\tex_lastnodechar:D</code>
1144	<code>__kernel_primitive:NN \lastnodesubtype</code>	<code>\tex_lastnodesubtype:D</code>
1145	<code>__kernel_primitive:NN \noautospace</code>	<code>\tex_noautospace:D</code>
1146	<code>__kernel_primitive:NN \noautoxspace</code>	<code>\tex_noautoxspace:D</code>
1147	<code>__kernel_primitive:NN \pagefistretch</code>	<code>\tex_pagefistretch:D</code>
1148	<code>__kernel_primitive:NN \postbreakpenalty</code>	<code>\tex_postbreakpenalty:D</code>
1149	<code>__kernel_primitive:NN \prebreakpenalty</code>	<code>\tex_prebreakpenalty:D</code>
1150	<code>__kernel_primitive:NN \ptexminorversion</code>	<code>\tex_ptexminorversion:D</code>
1151	<code>__kernel_primitive:NN \ptexrevision</code>	<code>\tex_ptexrevision:D</code>
1152	<code>__kernel_primitive:NN \ptexversion</code>	<code>\tex_ptexversion:D</code>
1153	<code>__kernel_primitive:NN \readpapersizespecial</code>	<code>\tex_readpapersizespecial:D</code>
1154	<code>__kernel_primitive:NN \scriptbaselineshiftfactor</code>	
1155	<code>\tex_scriptbaselineshiftfactor:D</code>	
1156	<code>__kernel_primitive:NN \scriptscriptbaselineshiftfactor</code>	
1157	<code>\tex_scriptscriptbaselineshiftfactor:D</code>	
1158	<code>__kernel_primitive:NN \showmode</code>	<code>\tex_showmode:D</code>
1159	<code>__kernel_primitive:NN \sjis</code>	<code>\tex_sjis:D</code>
1160	<code>__kernel_primitive:NN \tate</code>	<code>\tex_tate:D</code>
1161	<code>__kernel_primitive:NN \tbaselineshift</code>	<code>\tex_tbaselineshift:D</code>
1162	<code>__kernel_primitive:NN \textbaselineshiftfactor</code>	
1163	<code>\tex_textbaselineshiftfactor:D</code>	
1164	<code>__kernel_primitive:NN \tfont</code>	<code>\tex_tfont:D</code>
1165	<code>__kernel_primitive:NN \xkanjiskip</code>	<code>\tex_xkanjiskip:D</code>
1166	<code>__kernel_primitive:NN \xspcode</code>	<code>\tex_xspcode:D</code>
1167	<code>__kernel_primitive:NN \ybaselineshift</code>	<code>\tex_ybaselineshift:D</code>
1168	<code>__kernel_primitive:NN \yoko</code>	<code>\tex_yoko:D</code>
1169	<code>__kernel_primitive:NN \vfi</code>	<code>\tex_vfi:D</code>

Primitives from up $\mathbb{L}\mathbb{E}\mathbb{X}$.

1170	<code>__kernel_primitive:NN \currentcjktoken</code>	<code>\tex_currentcjktoken:D</code>
1171	<code>__kernel_primitive:NN \disablecjktoken</code>	<code>\tex_disablecjktoken:D</code>
1172	<code>__kernel_primitive:NN \enablecjktoken</code>	<code>\tex_enablecjktoken:D</code>
1173	<code>__kernel_primitive:NN \forcecjktoken</code>	<code>\tex_forcecjktoken:D</code>

```

1174 \__kernel_primitive:NN \kchar \tex_kchar:D
1175 \__kernel_primitive:NN \kchardef \tex_kchardef:D
1176 \__kernel_primitive:NN \kuten \tex_kuten:D
1177 \__kernel_primitive:NN \ucs \tex_ucs:D
1178 \__kernel_primitive:NN \uptexrevision \tex_uptexrevision:D
1179 \__kernel_primitive:NN \uptexversion \tex_uptexversion:D

```

Omega primitives provided by pTeX (listed separately mainly to allow understanding of their source).

```

1180 \__kernel_primitive:NN \odelcode \tex_odelcode:D
1181 \__kernel_primitive:NN \odelimiter \tex_odelimiter:D
1182 \__kernel_primitive:NN \omathaccent \tex_omathaccent:D
1183 \__kernel_primitive:NN \omathchar \tex_omathchar:D
1184 \__kernel_primitive:NN \omathchardef \tex_omathchardef:D
1185 \__kernel_primitive:NN \omathcode \tex_omathcode:D
1186 \__kernel_primitive:NN \oradical \tex_oradical:D

```

End of the “just the names” part of the source.

```

1187 </names | package>
1188 </names | tex>
1189 <*package>
1190 <*tex>

```

The job is done: close the group (using the primitive renamed!).

```

1191 \tex_endgroup:D

```

L^AT_EX 2_ε moves a few primitives, so these are sorted out. In newer versions of L^AT_EX 2_ε, expl3 is loaded rather early, so only some primitives are already renamed, so we need two tests here. At the beginning of the L^AT_EX 2_ε format, the primitives `\end` and `\input` are renamed, and only later on the other ones.

```

1192 \tex_ifdefined:D \@@end
1193 \tex_let:D \tex_end:D \@@end
1194 \tex_let:D \tex_input:D \@@input
1195 \tex_fi:D

```

If `\@@@hyph` is defined, we are loading expl3 in a pre-2020/10/01 release of L^AT_EX 2_ε, so a few other primitives have to be tested as well.

```

1196 \tex_ifdefined:D \@@hyph
1197 \tex_let:D \tex_everydisplay:D \frozen@everydisplay
1198 \tex_let:D \tex_everymath:D \frozen@everymath
1199 \tex_let:D \tex_hyphen:D \@@hyph
1200 \tex_let:D \tex_italiccorrection:D \@@italiccorr
1201 \tex_let:D \tex_underline:D \@@underline

```

The `\shipout` primitive is particularly tricky as a number of packages want to hook in here. First, we see if a sufficiently-new kernel has saved a copy: if it has, just use that. Otherwise, we need to check each of the possible packages/classes that might move it: here, we are looking for those which do *not* delay action to the `\AtBeginDocument` hook. (We cannot use `\primitive` as that doesn’t allow us to make a direct copy of the primitive *itself*.) As we know that L^AT_EX 2_ε is in use, we use its `\@tfor` loop here.

```

1202 \tex_ifdefined:D \@@shipout
1203 \tex_let:D \tex_shipout:D \@@shipout
1204 \tex_fi:D
1205 \tex_begingroup:D
1206 \tex_edef:D \l_tmpa_tl { \tex_string:D \shipout }

```

```

1207 \tex_edef:D \l_tmpb_tl { \tex_meaning:D \shipout }
1208 \tex_ifx:D \l_tmpa_tl \l_tmpb_tl
1209 \tex_else:D
1210 \tex_expandafter:D \@tfor \tex_expandafter:D \@tempa \tex_string:D :=
1211 \CROP@shipout
1212 \dup@shipout
1213 \GPTorg@shipout
1214 \LL@shipout
1215 \mem@oldshipout
1216 \open@shipout
1217 \pgfpages@originalshipout
1218 \pr@shipout
1219 \Shipout
1220 \verso@orig@shipout
1221 \do
1222 {
1223 \tex_edef:D \l_tmpb_tl
1224 { \tex_expandafter:D \tex_meaning:D \@tempa }
1225 \tex_ifx:D \l_tmpa_tl \l_tmpb_tl
1226 \tex_global:D \tex_expandafter:D \tex_let:D
1227 \tex_expandafter:D \tex_shipout:D \@tempa
1228 \tex_fi:D
1229 }
1230 \tex_fi:D
1231 \tex_endgroup:D

```

Some tidying up is needed for `\(pdf)tracingfonts`. Newer LuaTeX has this simply as `\tracingfonts`, but that is overwritten by the L^AT_EX 2_ε kernel. So any spurious definition has to be removed, then the real version saved either from the pdfTeX name or from LuaTeX. In the latter case, we leave `\@@tracingfonts` available; this might be useful and almost all L^AT_EX 2_ε users will have `expl3` loaded by `fontspec`. (We follow the usual kernel convention that `@@` is used for saved primitives.)

```

1232 \tex_let:D \tex_tracingfonts:D \tex_undefined:D
1233 \tex_ifdefined:D \pdftracingfonts
1234 \tex_let:D \tex_tracingfonts:D \pdftracingfonts
1235 \tex_else:D
1236 \tex_ifdefined:D \tex_directlua:D
1237 \tex_directlua:D { tex.enableprimitives("@@", {"tracingfonts"}) }
1238 \tex_let:D \tex_tracingfonts:D \@@tracingfonts
1239 \tex_fi:D
1240 \tex_fi:D
1241 \tex_fi:D

```

That is also true for the LuaTeX primitives under L^AT_EX 2_ε (depending on the format-building date). There are a few primitives that get the right names anyway so are missing here!

```

1242 \tex_ifdefined:D \luatexsuppressfontnotfounderror
1243 \tex_let:D \tex_alignmark:D \luatexalignmark
1244 \tex_let:D \tex_aligntab:D \luatexaligntab
1245 \tex_let:D \tex_attribute:D \luatexattribute
1246 \tex_let:D \tex_attributedef:D \luatexattributedef
1247 \tex_let:D \tex_catcodetable:D \luatexcatcodetable
1248 \tex_let:D \tex_clearmarks:D \luatexclearmarks
1249 \tex_let:D \tex_crampeddisplaystyle:D \luatexcrampeddisplaystyle

```



```

1250 \tex_let:D \tex_crampedscriptscriptstyle:D
1251 \luatexcrampedscriptscriptstyle
1252 \tex_let:D \tex_crampedscriptstyle:D \luatexcrampedscriptstyle
1253 \tex_let:D \tex_crampedtextstyle:D \luatexcrampedtextstyle
1254 \tex_let:D \tex_fontid:D \luatexfontid
1255 \tex_let:D \tex_formatname:D \luatexformatname
1256 \tex_let:D \tex_gleaders:D \luatexgleaders
1257 \tex_let:D \tex_initcatcodetable:D \luatexinitcatcodetable
1258 \tex_let:D \tex_latelua:D \luatexlatelua
1259 \tex_let:D \tex_luaescapestring:D \luatexluaescapestring
1260 \tex_let:D \tex_luafunction:D \luatexluafunction
1261 \tex_let:D \tex_mathstyle:D \luatexmathstyle
1262 \tex_let:D \tex_nokerns:D \luatexnokerns
1263 \tex_let:D \tex_noligs:D \luatexnoligs
1264 \tex_let:D \tex_outputbox:D \luatexoutputbox
1265 \tex_let:D \tex_pageleftoffset:D \luatexpageleftoffset
1266 \tex_let:D \tex_pagetopoffset:D \luatexpagetopoffset
1267 \tex_let:D \tex_postexhyphenchar:D \luatexpostexhyphenchar
1268 \tex_let:D \tex_posthyphenchar:D \luatexposthyphenchar
1269 \tex_let:D \tex_preexhyphenchar:D \luatexpreexhyphenchar
1270 \tex_let:D \tex_prehyphenchar:D \luatexprehyphenchar
1271 \tex_let:D \tex_savecatcodetable:D \luatexsavecatcodetable
1272 \tex_let:D \tex_scantextokens:D \luatexscantextokens
1273 \tex_let:D \tex_suppressifcsnameerror:D
1274 \luatexsuppressifcsnameerror
1275 \tex_let:D \tex_suppresslongerror:D \luatexsuppresslongerror
1276 \tex_let:D \tex_suppressmathparerror:D
1277 \luatexsuppressmathparerror
1278 \tex_let:D \tex_suppressoutererror:D \luatexsuppressoutererror
1279 \tex_let:D \tex_Uchar:D \luatexUchar
1280 \tex_let:D \tex_suppressfontnotfounderror:D
1281 \luatexsuppressfontnotfounderror

```

Which also covers those slightly odd ones.

```

1282 \tex_let:D \tex_bodydir:D \luatexbodydir
1283 \tex_let:D \tex_boxdir:D \luatexboxdir
1284 \tex_let:D \tex_leftghost:D \luatexleftghost
1285 \tex_let:D \tex_localbrokenpenalty:D \luatexlocalbrokenpenalty
1286 \tex_let:D \tex_localinterlinepenalty:D
1287 \luatexlocalinterlinepenalty
1288 \tex_let:D \tex_localleftbox:D \luatexlocalleftbox
1289 \tex_let:D \tex_localrightbox:D \luatexlocalrightbox
1290 \tex_let:D \tex_mathdir:D \luatexmathdir
1291 \tex_let:D \tex_pagebottomoffset:D \luatexpagebottomoffset
1292 \tex_let:D \tex_pagedir:D \luatexpagedir
1293 \tex_let:D \tex_pageheight:D \luatexpageheight
1294 \tex_let:D \tex_pagerightoffset:D \luatexpagerightoffset
1295 \tex_let:D \tex_pagewidth:D \luatexpagewidth
1296 \tex_let:D \tex_pardir:D \luatexpardir
1297 \tex_let:D \tex_rightghost:D \luatexrightghost
1298 \tex_let:D \tex_textdir:D \luatextextdir
1299 \tex_fi:D

```

Only pdfTeX and LuaTeX define \pdfmapfile and \pdfmapline: Tidy up the fact that some format-building processes leave a couple of questionable decisions about that!

```

1300 \tex_ifnum:D 0
1301 \tex_ifdefined:D \tex_pdftexversion:D 1 \tex_fi:D
1302 \tex_ifdefined:D \tex_luatexversion:D 1 \tex_fi:D
1303 = 0 %
1304 \tex_let:D \tex_mapfile:D \tex_undefined:D
1305 \tex_let:D \tex_mapline:D \tex_undefined:D
1306 \tex_fi:D

```

A few packages do unfortunate things to date-related primitives.

```

1307 \tex_begingroup:D
1308 \tex_edef:D \l_tmpa_tl { \tex_meaning:D \tex_time:D }
1309 \tex_edef:D \l_tmpb_tl { \tex_string:D \time }
1310 \tex_ifx:D \l_tmpa_tl \l_tmpb_tl
1311 \tex_else:D
1312 \tex_global:D \tex_let:D \tex_time:D \tex_undefined:D
1313 \tex_fi:D
1314 \tex_edef:D \l_tmpa_tl { \tex_meaning:D \tex_day:D }
1315 \tex_edef:D \l_tmpb_tl { \tex_string:D \day }
1316 \tex_ifx:D \l_tmpa_tl \l_tmpb_tl
1317 \tex_else:D
1318 \tex_global:D \tex_let:D \tex_day:D \tex_undefined:D
1319 \tex_fi:D
1320 \tex_edef:D \l_tmpa_tl { \tex_meaning:D \tex_month:D }
1321 \tex_edef:D \l_tmpb_tl { \tex_string:D \month }
1322 \tex_ifx:D \l_tmpa_tl \l_tmpb_tl
1323 \tex_else:D
1324 \tex_global:D \tex_let:D \tex_month:D \tex_undefined:D
1325 \tex_fi:D
1326 \tex_edef:D \l_tmpa_tl { \tex_meaning:D \tex_year:D }
1327 \tex_edef:D \l_tmpb_tl { \tex_string:D \year }
1328 \tex_ifx:D \l_tmpa_tl \l_tmpb_tl
1329 \tex_else:D
1330 \tex_global:D \tex_let:D \tex_year:D \tex_undefined:D
1331 \tex_fi:D
1332 \tex_endgroup:D

```

Up to v0.80, LuaTeX defines the pdfTeX version data: rather confusing. Removing them means that `\tex_pdftexversion:D` is a marker for pdfTeX alone: useful in engine-dependent code later.

```

1333 \tex_ifdefined:D \tex_luatexversion:D
1334 \tex_let:D \tex_pdftexbanner:D \tex_undefined:D
1335 \tex_let:D \tex_pdftexrevision:D \tex_undefined:D
1336 \tex_let:D \tex_pdftexversion:D \tex_undefined:D
1337 \tex_fi:D

```

`cslatex` moves a couple of primitives which we recover here; as there is no other marker, we can only work by looking for the names.

```

1338 \tex_ifdefined:D \orieveryjob
1339 \tex_let:D \tex_everyjob:D \orieveryjob
1340 \tex_fi:D
1341 \tex_ifdefined:D \oripdfoutput
1342 \tex_let:D \tex_pdfoutput:D \oripdfoutput
1343 \tex_fi:D

```

For ConTeXt, two tests are needed. Both Mark II and Mark IV move several primitives: these are all covered by the first test, again using `\end` as a marker. For Mark IV,

a few more primitives are moved: they are implemented using some Lua code in the current ConT_EXt.

```

1344 \tex_ifdefined:D \normalend
1345 \tex_let:D \tex_end:D \normalend
1346 \tex_let:D \tex_everyjob:D \normaleveryjob
1347 \tex_let:D \tex_input:D \normalinput
1348 \tex_let:D \tex_language:D \normallanguage
1349 \tex_let:D \tex_mathop:D \normalmathop
1350 \tex_let:D \tex_month:D \normalmonth
1351 \tex_let:D \tex_outer:D \normalouter
1352 \tex_let:D \tex_over:D \normalover
1353 \tex_let:D \tex_vcenter:D \normalvcenter
1354 \tex_let:D \tex_unexpanded:D \normalunexpanded
1355 \tex_let:D \tex_expanded:D \normalexpanded
1356 \tex_fi:D
1357 \tex_ifdefined:D \normalitaliccorrection
1358 \tex_let:D \tex_hoffset:D \normalhoffset
1359 \tex_let:D \tex_italiccorrection:D \normalitaliccorrection
1360 \tex_let:D \tex_voffset:D \normalvoffset
1361 \tex_let:D \tex_showtokens:D \normalshowtokens
1362 \tex_let:D \tex_bodydir:D \spac_directions_normal_body_dir
1363 \tex_let:D \tex_pagedir:D \spac_directions_normal_page_dir
1364 \tex_fi:D
1365 \tex_ifdefined:D \normalleft
1366 \tex_let:D \tex_left:D \normalleft
1367 \tex_let:D \tex_middle:D \normalmiddle
1368 \tex_let:D \tex_right:D \normalright
1369 \tex_fi:D
1370 </tex>

```

In LuaT_EX, we additionally emulate some primitives using Lua code.

```

1371 <lua>

```

`\tex_strcmp:D` Compare two strings, expanding to 0 if they are equal, -1 if the first one is smaller and 1 if the second one is smaller. Here “smaller” refers to codepoint order which does not correspond to the user expected order for most non-ASCII strings.

```

1372 local minus_tok = token.new(string.byte'-' , 12)
1373 local zero_tok = token.new(string.byte'0' , 12)
1374 local one_tok = token.new(string.byte'1' , 12)
1375 luacmd('tex_strcmp:D', function()
1376   local first = scan_string()
1377   local second = scan_string()
1378   if first < second then
1379     put_next(minus_tok, one_tok)
1380   else
1381     put_next(first == second and zero_tok or one_tok)
1382   end
1383 end, 'global')

```

(End definition for `\tex_strcmp:D`. This function is documented on page ??.)

`\tex_Ucharcat:D` Creating arbitrary chars using `tex.cprint`. The alternative approach using `token.put_next(token.cre` would be about 10% slower.

```

1384 local cprint = tex.cprint
1385 luacmd('tex_Ucharcat:D', function()
1386     local charcode = scan_int()
1387     local catcode = scan_int()
1388     cprint(catcode, utf8_char(charcode))
1389 end, 'global')

```

(End definition for \tex_Ucharcat:D. This function is documented on page ??.)

\tex_filesize:D Wrap the function from ltxutils.

```

1390 luacmd('tex_filesize:D', function()
1391     local size = filesize(scan_string())
1392     if size then write(size) end
1393 end, 'global')

```

(End definition for \tex_filesize:D. This function is documented on page ??.)

\tex_md5sum:D There are two cases: Either hash a file or a string. Both are already implemented in l3luatex or built-in.

```

1394 luacmd('tex_md5sum:D', function()
1395     local hash
1396     if scan_keyword"file" then
1397         hash = filemd5sum(scan_string())
1398     else
1399         hash = md5_HEX(scan_string())
1400     end
1401     if hash then write(hash) end
1402 end, 'global')

```

(End definition for \tex_md5sum:D. This function is documented on page ??.)

\tex_filemoddate:D A primitive for getting the modification date of a file.

```

1403 luacmd('tex_filemoddate:D', function()
1404     local date = filemoddate(scan_string())
1405     if date then write(date) end
1406 end, 'global')

```

(End definition for \tex_filemoddate:D. This function is documented on page ??.)

\tex_filedump:D An emulated primitive for getting a hexdump from a (partial) file. The length has a default of 0. This is consistent with pdfTeX, but it effectively makes the primitive useless without an explicit `length`. Therefore we allow the keyword `whole` to be used instead of a length, indicating that the whole remaining file should be read.

```

1407 luacmd('tex_filedump:D', function()
1408     local offset = scan_keyword'offset' and scan_int() or nil
1409     local length = scan_keyword'length' and scan_int()
1410                 or not scan_keyword'whole' and 0 or nil
1411     local data = filedump(scan_string(), offset, length)
1412     if data then write(data) end
1413 end, 'global')

```

(End definition for \tex_filedump:D. This function is documented on page ??.)

```

1414 </lua>
1415 </package>

```

3 Internal kernel functions

<hr/> <code>_kernel_chk_cs_exist:N</code> <hr/> <code>_kernel_chk_cs_exist:c</code> <hr/>	<code>_kernel_chk_cs_exist:N <cs></code> This function is only created if debugging is enabled. It checks that <code><cs></code> exists according to the criteria for <code>\cs_if_exist_p:N</code> , and if not raises a kernel-level error.
<hr/> <code>_kernel_chk_defined:NT</code> <hr/>	<code>_kernel_chk_defined:NT <variable> {\true code}</code> If <code><variable></code> is not defined (according to <code>\cs_if_exist:NTF</code>), this triggers an error, otherwise the <code><true code></code> is run.
<hr/> <code>_kernel_chk_expr:nNnN</code> <hr/>	<code>_kernel_chk_expr:nNnN {\<expr>} <eval> {\<convert>} <caller></code> This function is only created if debugging is enabled. By default it is equivalent to <code>\use_i:nnnn</code> . When expression checking is enabled, it leaves in the input stream the result of <code>\tex_the:D <eval> <expr> \tex_relax:D</code> after checking that no token was left over. If any token was not taken as part of the expression, there is an error message displaying the result of the evaluation as well as the <code><caller></code> . For instance <code><eval></code> can be <code>_int_eval:w</code> and <code><caller></code> can be <code>\int_eval:n</code> or <code>\int_set:Nn</code> . The argument <code><convert></code> is empty except for mu expressions where it is <code>\tex_mutoglu:D</code> , used for internal purposes.
<hr/> <code>_kernel_cs_parm_from_arg_count:nnF</code> <hr/>	<code>_kernel_cs_parm_from_arg_count:nnF {\<follow-on>} {\<args>} {\<>false code>}</code> Evaluates the number of <code><args></code> and leaves the <code><follow-on></code> code followed by a brace group containing the required number of primitive parameter markers (<code>#1</code> , etc.). If the number of <code><args></code> is outside the range <code>[0,9]</code> , the <code><>false code></code> is inserted <i>instead</i> of the <code><follow-on></code> .
<hr/> <code>_kernel_dependency_version_check:Nn</code> <hr/> <code>_kernel_dependency_version_check:nn</code> <hr/>	<code>_kernel_dependency_version_check:Nn {\<date>} {\<file>}</code> <code>_kernel_dependency_version_check:nn {\<date>} {\<file>}</code> Checks if the loaded version of the <code>expl3</code> kernel is at least <code><date></code> , required by <code><file></code> . If the kernel date is older than <code><date></code> , the loading of <code><file></code> is aborted and an error is raised.
<hr/> <code>_kernel_deprecation_code:nn</code> <hr/>	<code>_kernel_deprecation_code:nn {\<error code>} {\<working code>}</code> Stores both an <code><error></code> and <code><working></code> definition for given material such that they can be exchanged by <code>\debug_on:</code> and <code>\debug_off:</code> .
<hr/> <code>_kernel_exp_not:w *</code> <hr/>	<code>_kernel_exp_not:w <expandable tokens> {\<content>}</code> Carries out expansion on the <code><expandable tokens></code> before preventing further expansion of the <code><content></code> as for <code>\exp_not:n</code> . Typically, the <code><expandable tokens></code> will alter the nature of the <code><content></code> , <i>i.e.</i> allow it to be generated in some way.
<code>\l__kernel_expl_bool</code>	A boolean which records the current code syntax status: <code>true</code> if currently inside a code environment. This variable should only be set by <code>\ExplSyntaxOn/\ExplSyntaxOff</code> . (End definition for <code>\l__kernel_expl_bool</code> .)
<code>\c__kernel_expl_date_tl</code>	A token list containing the release date of the <code>l3kernel</code> preloaded in $\text{\LaTeX 2}_{\epsilon}$ used to check if dependencies match.

(End definition for `\c__kernel_expl_date_tl.`)

`__kernel_file_missing:n`

`__kernel_file_missing:n` $\{\langle name \rangle\}$

Expands the $\langle name \rangle$ as per `__kernel_file_name_sanitize:nN` then produces an error message indicating that this file was not found.

`__kernel_file_name_sanitize:nN`

`__kernel_file_name_sanitize:nN` $\{\langle name \rangle\}$ $\langle str var \rangle$

For converting a $\langle name \rangle$ to a string where active characters are treated as strings.

`__kernel_file_input_push:n`

`__kernel_file_input_push:n` $\{\langle name \rangle\}$

`__kernel_file_input_pop:`

`__kernel_file_input_pop:`

Used to push and pop data from the internal file stack: needed only in package mode, where interfacing with the L^AT_EX 2_ε kernel is necessary.

`__kernel_int_add:nnn` \star

`__kernel_int_add:nnn` $\{\langle integer_1 \rangle\}$ $\{\langle integer_2 \rangle\}$ $\{\langle integer_3 \rangle\}$

Expands to the result of adding the three $\langle integers \rangle$ (which must be suitable input for `\int_eval:w`), avoiding intermediate overflow. Overflow occurs only if the overall result is outside $[-2^{31}+1, 2^{31}-1]$. The $\langle integers \rangle$ may be of the form `\int_eval:w ... \scan_stop:` but may be evaluated more than once.

`__kernel_intarray_gset:Nnn`

`__kernel_intarray_gset:Nnn` $\langle intarray var \rangle$ $\{\langle index \rangle\}$ $\{\langle value \rangle\}$

New: 2018-03-31

Faster version of `\intarray_gset:Nnn`. Stores the $\langle value \rangle$ into the $\langle integer array variable \rangle$ at the $\langle position \rangle$. The $\langle index \rangle$ and $\langle value \rangle$ must be suitable for a direct assignment to a T_EX count register, for instance expanding to an integer denotation or obtained through the primitive `\numexpr` (which may be un-terminated). No bound checking is performed: the caller is responsible for ensuring that the $\langle position \rangle$ is between 1 and the `\intarray_count:N`, and the $\langle value \rangle$'s absolute value is at most $2^{30}-1$. Assignments are always global.

`__kernel_intarray_item:Nn` \star

`__kernel_intarray_item:Nn` $\langle intarray var \rangle$ $\{\langle index \rangle\}$

New: 2018-03-31

Faster version of `\intarray_item:Nn`. Expands to the integer entry stored at the $\langle index \rangle$ in the $\langle integer array variable \rangle$. The $\langle index \rangle$ must be suitable for a direct assignment to a T_EX count register and must be between 1 and the `\intarray_count:N`, lest a low-level T_EX error occur.

`__kernel_intarray_range_to_clist:Nnn` \star

`__kernel_intarray_range_to_clist:Nnn` $\langle intarray var \rangle$ $\{\langle start index \rangle\}$ $\{\langle end index \rangle\}$

New: 2020-07-12

Converts to integer denotations separated by commas the entries of the $\langle intarray \rangle$ from positions $\langle start index \rangle$ to $\langle end index \rangle$ included. The $\langle start index \rangle$ and $\langle end index \rangle$ must be suitable for a direct assignment to a T_EX count register, must be between 1 and the `\intarray_count:N`, and be suitably ordered. All tokens have category code other.

`__kernel_intarray_gset_range_from_clist:Nnn`

New: 2020-07-12

`__kernel_intarray_gset_range_from_clist:Nnn`
`<intarray var> {<start index>} {<integer clist>}`

Stores the entries of the `<clist>` as entries of the `<intarray var>` starting from the `<start index>`, upwards. This is done without any bound checking. The `<start index>` and all entries of the `<integer comma list>` (which do not undergo space trimming and brace stripping as in normal clist mappings) must be suitable for a direct assignment to a T_EX count register. An empty entry may stop the loop.

`__kernel_ior_open:Nn`
`__kernel_ior_open:No`

`__kernel_ior_open:Nn <stream> {<file name>}`

This function has identical syntax to the public version. However, it does not take precautions against active characters in the `<file name>`, and it does not attempt to add a `<path>` to the `<file name>`: it is therefore intended to be used by higher-level functions which have already fully expanded the `<file name>` and which need to perform multiple open or close operations. See for example the implementation of `\file_get_full_name:nN`,

`__kernel_iow_with:Nnn`

`__kernel_iow_with:Nnn <integer> {<value>} {<code>}`

If the `<integer>` is equal to the `<value>` then this function simply runs the `<code>`. Otherwise it saves the current value of the `<integer>`, sets it to the `<value>`, runs the `<code>`, and restores the `<integer>` to its former value. This is used to ensure that the `\newlinechar` is 10 when writing to a stream, which lets `\iow_newline:` work, and that `\errorcontextlines` is -1 when displaying a message.

`__kernel_msg_new:nnnn`
`__kernel_msg_new:nnn`

`__kernel_msg_new:nnnn {<module>} {<message>} {<text>} {<more text>}`

Creates a kernel `<message>` for a given `<module>`. The message is defined to first give `<text>` and then `<more text>` if the user requests it. If no `<more text>` is available then a standard text is given instead. Within `<text>` and `<more text>` four parameters (#1 to #4) can be used: these will be supplied and expanded at the time the message is used. An error is raised if the `<message>` already exists.

`__kernel_msg_set:nnnn`
`__kernel_msg_set:nnn`

`__kernel_msg_set:nnnn {<module>} {<message>} {<text>} {<more text>}`

Sets up the text for a kernel `<message>` for a given `<module>`. The message is defined to first give `<text>` and then `<more text>` if the user requests it. If no `<more text>` is available then a standard text is given instead. Within `<text>` and `<more text>` four parameters (#1 to #4) can be used: these will be supplied and expanded at the time the message is used.

`__kernel_msg_fatal:nnnnnn`
`__kernel_msg_fatal:nnxxxx`
`__kernel_msg_fatal:nnnnn`
`__kernel_msg_fatal:nnxxx`
`__kernel_msg_fatal:nnnn`
`__kernel_msg_fatal:nnxx`
`__kernel_msg_fatal:nnn`
`__kernel_msg_fatal:nnx`
`__kernel_msg_fatal:nn`

`__kernel_msg_fatal:nnnnnn {<module>} {<message>} {<arg one>} {<arg two>} {<arg three>} {<arg four>}`

Issues kernel `<module>` error `<message>`, passing `<arg one>` to `<arg four>` to the text-creating functions. After issuing a fatal error the T_EX run halts. Cannot be redirected.

```

\_kernel_msg_critical:nnnnnn \_kernel_msg_critical:nnnnnn {\module}} {\message}} {\arg one}} {\arg
\_kernel_msg_critical:nnxxxx two}} {\arg three}} {\arg four}}
\_kernel_msg_critical:nnnnn
\_kernel_msg_critical:nnxxx
\_kernel_msg_critical:nnnn
\_kernel_msg_critical:nnxx
\_kernel_msg_critical:nnn
\_kernel_msg_critical:nnx
\_kernel_msg_critical:nn

```

Issues kernel $\langle module \rangle$ error $\langle message \rangle$, passing $\langle arg one \rangle$ to $\langle arg four \rangle$ to the text-creating functions. After issuing a critical error, T_EX stops reading the current input file. Cannot be redirected.

```

\_kernel_msg_error:nnnnnn \_kernel_msg_error:nnnnnn {\module}} {\message}} {\arg one}} {\arg two}} {\arg
\_kernel_msg_error:nnxxxx three}} {\arg four}}
\_kernel_msg_error:nnnnn
\_kernel_msg_error:nnxxx
\_kernel_msg_error:nnnn
\_kernel_msg_error:nnxx
\_kernel_msg_error:nnn
\_kernel_msg_error:nnx
\_kernel_msg_error:nn

```

Issues kernel $\langle module \rangle$ error $\langle message \rangle$, passing $\langle arg one \rangle$ to $\langle arg four \rangle$ to the text-creating functions. The error stops processing and issues the text at the terminal. After user input, the run continues. Cannot be redirected.

```

\_kernel_msg_warning:nnnnnn \_kernel_msg_warning:nnnnnn {\module}} {\message}} {\arg one}} {\arg
\_kernel_msg_warning:nnxxxx two}} {\arg three}} {\arg four}}
\_kernel_msg_warning:nnnnn
\_kernel_msg_warning:nnxxx
\_kernel_msg_warning:nnnn
\_kernel_msg_warning:nnxx
\_kernel_msg_warning:nnn
\_kernel_msg_warning:nnx
\_kernel_msg_warning:nn

```

Issues kernel $\langle module \rangle$ warning $\langle message \rangle$, passing $\langle arg one \rangle$ to $\langle arg four \rangle$ to the text-creating functions. The warning text is added to the log file, but the T_EX run is not interrupted.

```

\_kernel_msg_info:nnnnnn \_kernel_msg_info:nnnnnn {\module}} {\message}} {\arg one}} {\arg two}} {\arg
\_kernel_msg_info:nnxxxx three}} {\arg four}}
\_kernel_msg_info:nnnnn
\_kernel_msg_info:nnxxx
\_kernel_msg_info:nnnn
\_kernel_msg_info:nnxx
\_kernel_msg_info:nnn
\_kernel_msg_info:nnx
\_kernel_msg_info:nn

```

Issues kernel $\langle module \rangle$ information $\langle message \rangle$, passing $\langle arg one \rangle$ to $\langle arg four \rangle$ to the text-creating functions. The information text is added to the log file.

```

\__kernel_msg_expandable_error:nnnnnn * \__kernel_msg_expandable_error:nnnnnn {\module} {\message}
\__kernel_msg_expandable_error:nnffff * {\arg one} {\arg two} {\arg three} {\arg four}
\__kernel_msg_expandable_error:nnnnn *
\__kernel_msg_expandable_error:nnfff *
\__kernel_msg_expandable_error:nnnn *
\__kernel_msg_expandable_error:nnff *
\__kernel_msg_expandable_error:nnn *
\__kernel_msg_expandable_error:nnf *
\__kernel_msg_expandable_error:nn *

```

Issues an error, passing *⟨arg one⟩* to *⟨arg four⟩* to the text-creating functions. The resulting string must be much shorter than a line, otherwise it is cropped.

`\g__kernel_prg_map_int` This integer is used by non-expandable mapping functions to track the level of nesting in force. The functions `\⟨type⟩_map_1:w`, `\⟨type⟩_map_2:w`, *etc.*, labelled by `\g__kernel_prg_map_int` hold functions to be mapped over various list datatypes in inline and variable mappings.

(End definition for `\g__kernel_prg_map_int`.)

`__kernel_quark_new_test:N`

`__kernel_quark_new_test:N \<name>:<arg spec>`

Defines a quark-test function `\<name>:<arg spec>` which tests if its argument is `\q__<namespace>_recursion_tail`, then acts accordingly, as described below for each possible `<arg spec>`.

The `<namespace>` is determined as the first (nonempty) `_`-delimited word in `<name>` and is used internally in the definition of auxiliaries. The function `__kernel_quark_new_test:N` does *not* define the `\q__<namespace>_recursion_tail` and `\q__<namespace>_recursion_stop` quarks. They should be manually defined with `\quark_new:N`.

There are 6 different types of quark-test functions. Which one is defined depends on the `<arg spec>`, which *must* be one of the options listed now. Four of them are modeled after `\quark_if_recursion_tail:(N|n)` and `\quark_if_recursion_tail_do:(N|n)n`.

`n` defines `\<name>:n` such that it checks if #1 contains only `\q__<namespace>_recursion_tail`, and if so consumes all tokens up to `\q__<namespace>_recursion_stop` (c.f. `\quark_if_recursion_tail_stop:n`).

`nn` defines `\<name>:nn` such that it checks if #1 contains only `\q__<namespace>_recursion_tail`, and if so consumes all tokens up to `\q__<namespace>_recursion_stop`, then executes the code #2 after that (c.f. `\quark_if_recursion_tail_stop_do:nn`).

`N` defines `\<name>:N` such that it checks if #1 is `\q__<namespace>_recursion_tail`, and if so consumes all tokens up to `\q__<namespace>_recursion_stop` (c.f. `\quark_if_recursion_tail_stop:N`).

`Nn` defines `\<name>:Nn` such that it checks if #1 is `\q__<namespace>_recursion_tail`, and if so consumes all tokens up to `\q__<namespace>_recursion_stop`, then executes the code #2 after that (c.f. `\quark_if_recursion_tail_stop_do:Nn`).

The last two are modeled after `\quark_if_recursion_tail_break:(n|N)N`, and in those cases the quark `\q__<namespace>_recursion_stop` is not used (and thus needs not be defined).

`nN` defines `\<name>:nN` such that it checks if #1 contains only `\q__<namespace>_recursion_tail`, and if so uses the `\<type>_map_break:` function #2.

`NN` defines `\<name>:NN` such that it checks if #1 is `\q__<namespace>_recursion_tail`, and if so uses the `\<type>_map_break:` function #2.

Any other signature, as well as a function without signature are errors, and in such case the definition is aborted.

<u><code>__kernel_quark_new_conditional:Nn</code></u>	<code>__kernel_quark_new_conditional:Nn</code> <code>__<namespace>_quark_if_<name>:<arg spec> {<conditions>}</code>
--	--

Defines a collection of quark conditionals that test if their argument is the quark `\q_`
`__<namespace>_<name>` and perform suitable actions. The `<conditions>` are a comma-separated list of one or more of p, T, F, and TF, and one conditional is defined for each `<condition>` in the list, as described for `\prg_new_conditional:Npnn`. The conditionals are defined using `\prg_new_conditional:Npnn`, so that their name is obtained by adding p, T, F, or TF to the base name `__<namespace>_quark_if_<name>:<arg spec>`.

The first argument of `__kernel_quark_new_conditional:Nn` must contain `_quark_if_` and `:`, as these markers are used to determine the `<name>` of the quark `\q_`
`__<namespace>_<name>` to be tested. This quark should be manually defined with `\quark_new:N`, as `__kernel_quark_new_conditional:Nn` does *not* define it.

The function `__kernel_quark_new_conditional:Nn` can define 2 different types of quark conditionals. Which one is defined depends on the `<arg spec>`, which *must* be one of the following options, modeled after `\quark_if_nil:(N|n)(TF)`.

`n` defines `__<namespace>_quark_if_<name>:n(TF)` such that it checks if #1 contains only `\q_`
`__<namespace>_<name>`, and executes the proper conditional branch.

`N` defines `__<namespace>_quark_if_<name>:N(TF)` such that it checks if #1 is `\q_`
`__<namespace>_<name>`, and executes the proper conditional branch.

Any other signature, as well as a function without signature are errors, and in such case the definition is aborted.

`\c__kernel_randint_max_int` Maximal allowed argument to `__kernel_randint:n`. Equal to $2^{17} - 1$.

(End definition for `\c__kernel_randint_max_int`.)

<u><code>__kernel_randint:n</code></u>	<code>__kernel_randint:n {<max>}</code>
---	--

Used in an integer expression this gives a pseudo-random number between 1 and `<max>` included. One must have $\langle max \rangle \leq 2^{17} - 1$. The `<max>` must be suitable for `\int_value:w` (and any `\int_eval:w` must be terminated by `\scan_stop:` or equivalent).

<u><code>__kernel_randint:nn</code></u>	<code>__kernel_randint:nn {<min>} {<max>}</code>
--	---

Used in an integer expression this gives a pseudo-random number between `<min>` and `<max>` included. The `<min>` and `<max>` must be suitable for `\int_value:w` (and any `\int_eval:w` must be terminated by `\scan_stop:` or equivalent). For small ranges $R = \langle max \rangle - \langle min \rangle + 1 \leq 2^{17} - 1$, `<min> - 1 + __kernel_randint:n{R}` is faster.

<u><code>__kernel_register_show:N</code></u> <u><code>__kernel_register_show:c</code></u>	<code>__kernel_register_show:N <register></code>
--	---

Used to show the contents of a T_EX register at the terminal, formatted such that internal parts of the mechanism are not visible.

<u><code>__kernel_register_log:N</code></u> <u><code>__kernel_register_log:c</code></u>	<code>__kernel_register_log:N <register></code>
--	--

Used to write the contents of a T_EX register to the log file in a form similar to `__kernel_register_show:N`.

<code>_kernel_str_to_other:n</code> ★	<code>_kernel_str_to_other:n {\token list}</code>
--	--

Converts the $\langle token list \rangle$ to a $\langle other string \rangle$, where spaces have category code “other”. This function can be f-expanded without fear of losing a leading space, since spaces do not have category code 10 in its result. It takes a time quadratic in the character count of the string.

<code>_kernel_str_to_other_fast:n</code> ☆	<code>_kernel_str_to_other_fast:n {\token list}</code>
---	---

Same behaviour `_kernel_str_to_other:n` but only restricted-expandable. It takes a time linear in the character count of the string.

<code>_kernel_tl_to_str:w</code> ★	<code>_kernel_tl_to_str:w \expandable tokens {\tokens}</code>
-------------------------------------	--

Carries out expansion on the $\langle expandable tokens \rangle$ before conversion of the $\langle tokens \rangle$ to a string as describe for `\tl_to_str:n`. Typically, the $\langle expandable tokens \rangle$ will alter the nature of the $\langle tokens \rangle$, *i.e.* allow it to be generated in some way. This function requires only a single expansion.

<code>_kernel_tl_set:Nx</code> <code>_kernel_tl_gset:Nx</code>	<code>_kernel_tl_set:Nx \tl var {\tokens}</code>
---	---

Fully expands $\langle tokens \rangle$ and assigns the result to $\langle tl var \rangle$. $\langle tokens \rangle$ must be given in braces and there must be no token between $\langle tl var \rangle$ and $\langle tokens \rangle$.

4 Kernel backend functions

These functions are required to pass information to the backend. The nature of these means that they are defined only when the relevant backend is in use.

<code>_kernel_backend_literal:n</code> <code>_kernel_backend_literal:(e x)</code>	<code>_kernel_backend_literal:n {\content}</code>
--	--

Adds the $\langle content \rangle$ literally to the current vertical list as a whatsit. The nature of the $\langle content \rangle$ will depend on the backend in use.

<code>_kernel_backend_literal_postscript:n</code> <code>_kernel_backend_literal_postscript:x</code>	<code>_kernel_backend_literal_postscript:n {\PostScript}</code>
--	--

Adds the $\langle PostScript \rangle$ literally to the current vertical list as a whatsit. No positioning is applied.

<code>_kernel_backend_literal_pdf:n</code> <code>_kernel_backend_literal_pdf:x</code>	<code>_kernel_backend_literal_pdf:n {\PDF instructions}</code>
--	---

Adds the $\langle PDF instructions \rangle$ literally to the current vertical list as a whatsit. No positioning is applied.

<code>_kernel_backend_literal_svg:n</code> <code>_kernel_backend_literal_svg:x</code>	<code>_kernel_backend_literal_svg:n {\SVG instructions}</code>
--	---

Adds the $\langle SVG instructions \rangle$ literally to the current vertical list as a whatsit. No positioning is applied.

<code>__kernel_backend_postscript:n</code>	<code>__kernel_backend_postscript:n {⟨PostScript⟩}</code>
<code>__kernel_backend_postscript:x</code>	

Adds the $\langle PostScript \rangle$ to the current vertical list as a whatsit. The PostScript reference point is adjusted to match the current position. The PostScript is inserted inside a SDict begin/end pair.

<code>__kernel_backend_align_begin:</code>	<code>__kernel_backend_align_begin:</code>
<code>__kernel_backend_align_end:</code>	<code>⟨PostScript literals⟩</code>
	<code>__kernel_backend_align_end:</code>

Arranges to align the PostScript and DVI current positions and scales.

<code>__kernel_backend_scope_begin:</code>	<code>__kernel_backend_scope_begin:</code>
<code>__kernel_backend_scope_end:</code>	<code>⟨content⟩</code>
	<code>__kernel_backend_scope_end:</code>

Creates a scope for instructions at the backend level.

<code>__kernel_backend_matrix:n</code>	<code>__kernel_backend_matrix:n {⟨matrix⟩}</code>
<code>__kernel_backend_matrix:x</code>	

Applies the $\langle matrix \rangle$ to the current transformation matrix.

<code>\g__kernel_backend_header_bool</code>

Specifies whether to write headers for the backend.

<code>\l__kernel_color_stack_int</code>	The color stack used in pdfTeX and LuaTeX for the main color.
---	---

5 l3basics implementation

¹⁴¹⁶ $\langle *package \rangle$

5.1 Renaming some TeX primitives (again)

Having given all the TeX primitives a consistent name, we need to give sensible names to the ones we actually want to use. These will be defined as needed in the appropriate modules, but we do a few now, just to get started.⁷

<code>\if_true:</code>	Then some conditionals.	
<code>\if_false:</code>	¹⁴¹⁷ <code>\tex_let:D \if_true:</code>	<code>\tex_iftrue:D</code>
<code>\or:</code>	¹⁴¹⁸ <code>\tex_let:D \if_false:</code>	<code>\tex_iffalse:D</code>
<code>\else:</code>	¹⁴¹⁹ <code>\tex_let:D \or:</code>	<code>\tex_or:D</code>
<code>\fi:</code>	¹⁴²⁰ <code>\tex_let:D \else:</code>	<code>\tex_else:D</code>
<code>\reverse_if:N</code>	¹⁴²¹ <code>\tex_let:D \fi:</code>	<code>\tex_fi:D</code>
<code>\if:w</code>	¹⁴²² <code>\tex_let:D \reverse_if:N</code>	<code>\tex_unless:D</code>
<code>\if_charcode:w</code>	¹⁴²³ <code>\tex_let:D \if:w</code>	<code>\tex_if:D</code>
<code>\if_catcode:w</code>	¹⁴²⁴ <code>\tex_let:D \if_charcode:w</code>	<code>\tex_if:D</code>
<code>\if_meaning:w</code>	¹⁴²⁵ <code>\tex_let:D \if_catcode:w</code>	<code>\tex_ifcat:D</code>
	¹⁴²⁶ <code>\tex_let:D \if_meaning:w</code>	<code>\tex_ifx:D</code>
	¹⁴²⁷ <code>\tex_let:D \if_bool:N</code>	<code>\tex_ifodd:D</code>

⁷This renaming gets expensive in terms of csname usage, an alternative scheme would be to just use the `\tex_...:D` name in the cases where no good alternative exists.

(End definition for `\if_true:` and others. These functions are documented on page 23.)

```

\if_mode_math:  TEX lets us detect some if its modes.
\if_mode_horizontal:  \tex_let:D \if_mode_math:      \tex_ifmmode:D
\if_mode_vertical:    \tex_let:D \if_mode_horizontal: \tex_ifhmode:D
\if_mode_inner:       \tex_let:D \if_mode_vertical:   \tex_ifvmode:D
                    1431 \tex_let:D \if_mode_inner:      \tex_ifinner:D

```

(End definition for `\if_mode_math:` and others. These functions are documented on page 23.)

```

\if_cs_exist:N    Building csnames and testing if control sequences exist.
\if_cs_exist:w    1432 \tex_let:D \if_cs_exist:N      \tex_ifdefined:D
                  \cs:w  \tex_let:D \if_cs_exist:w      \tex_ifcsname:D
                  \cs_end: \tex_let:D \cs:w            \tex_csname:D
                  1435 \tex_let:D \cs_end:            \tex_endcsname:D

```

(End definition for `\if_cs_exist:N` and others. These functions are documented on page 23.)

```

\exp_after:wN    The five \exp_ functions are used in the l3expan module where they are described.
\exp_not:N       1436 \tex_let:D \exp_after:wN      \tex_expandafter:D
\exp_not:n       1437 \tex_let:D \exp_not:N        \tex_noexpand:D
                  1438 \tex_let:D \exp_not:n        \tex_unexpanded:D
                  1439 \tex_let:D \exp:w          \tex_romannumeral:D
                  1440 \tex_chardef:D \exp_end: = 0 ~

```

(End definition for `\exp_after:wN`, `\exp_not:N`, and `\exp_not:n`. These functions are documented on page 33.)

```

\token_to_meaning:N Examining a control sequence or token.
\cs_meaning:N       1441 \tex_let:D \token_to_meaning:N \tex_meaning:D
                  1442 \tex_let:D \cs_meaning:N      \tex_meaning:D

```

(End definition for `\token_to_meaning:N` and `\cs_meaning:N`. These functions are documented on page 134.)

```

\tl_to_str:n      Making strings.
\token_to_str:N   1443 \tex_let:D \tl_to_str:n      \tex_detokenize:D
\__kernel_tl_to_str:w 1444 \tex_let:D \token_to_str:N      \tex_string:D
                  1445 \tex_let:D \__kernel_tl_to_str:w \tex_detokenize:D

```

(End definition for `\tl_to_str:n`, `\token_to_str:N`, and `__kernel_tl_to_str:w`. These functions are documented on page 51.)

```

\scan_stop:       The next three are basic functions for which there also exist versions that are safe inside
\group_begin:     alignments. These safe versions are defined in the l3prg module.
\group_end:       1446 \tex_let:D \scan_stop:      \tex_relax:D
                  1447 \tex_let:D \group_begin:    \tex_begingroup:D
                  1448 \tex_let:D \group_end:      \tex_endgroup:D

```

(End definition for `\scan_stop:`, `\group_begin:`, and `\group_end:`. These functions are documented on page 9.)

```
1449 <@@=int>
```

```

\if_int_compare:w For integers.
\__int_to_roman:w 1450 \tex_let:D \if_int_compare:w \tex_ifnum:D
                  1451 \tex_let:D \__int_to_roman:w \tex_romannumeral:D

```

(End definition for `\if_int_compare:w` and `__int_to_roman:w`. This function is documented on page 101.)

`\group_insert_after:N` Adding material after the end of a group.

```
1452 \tex_let:D \group_insert_after:N \tex_aftergroup:D
```

(End definition for `\group_insert_after:N`. This function is documented on page 9.)

`\exp_args:Nc` Discussed in l3expan, but needed much earlier.

`\exp_args:cc`

```
1453 \tex_long:D \tex_def:D \exp_args:Nc #1#2
1454 { \exp_after:wN #1 \cs:w #2 \cs_end: }
1455 \tex_long:D \tex_def:D \exp_args:cc #1#2
1456 { \cs:w #1 \exp_after:wN \cs_end: \cs:w #2 \cs_end: }
```

(End definition for `\exp_args:Nc` and `\exp_args:cc`. These functions are documented on page 29.)

`\token_to_meaning:c` A small number of variants defined by hand. Some of the necessary functions (`\use_i:nn`, `\use_ii:nn`, and `\exp_args:Nnc`) are not defined at that point yet, but will be defined before those variants are used. The `\cs_meaning:c` command must check for an undefined control sequence to avoid defining it mistakenly.

`\token_to_str:c`
`\cs_meaning:c`

```
1457 \tex_def:D \token_to_str:c { \exp_args:Nc \token_to_str:N }
1458 \tex_long:D \tex_def:D \cs_meaning:c #1
1459 {
1460   \if_cs_exist:w #1 \cs_end:
1461     \exp_after:wN \use_i:nn
1462   \else:
1463     \exp_after:wN \use_ii:nn
1464   \fi:
1465   { \exp_args:Nc \cs_meaning:N {#1} }
1466   { \tl_to_str:n {undefined} }
1467 }
1468 \tex_let:D \token_to_meaning:c = \cs_meaning:c
```

(End definition for `\token_to_meaning:N`. This function is documented on page 134.)

5.2 Defining some constants

`\c_zero_int` We need the constant `\c_zero_int` which is used by some functions in the l3alloc module. The rest are defined in the l3int module – at least for the ones that can be defined with `\tex_chardef:D` or `\tex_mathchardef:D`. For other constants the l3int module is required but it can't be used until the allocation has been set up properly!

```
1469 \tex_chardef:D \c_zero_int = 0 ~
```

(End definition for `\c_zero_int`. This variable is documented on page 100.)

`\c_max_register_int` This is here as this particular integer is needed both in package mode and to bootstrap l3alloc, and is documented in l3int. LuaTeX and those which contain parts of the Omega extensions have more registers available than ε -TeX.

```
1470 \tex_ifdefined:D \tex_luatexversion:D
1471 \tex_chardef:D \c_max_register_int = 65 535 ~
1472 \tex_else:D
1473 \tex_ifdefined:D \tex_omathchardef:D
1474 \tex_omathchardef:D \c_max_register_int = 65535 ~
1475 \tex_else:D
```

```

1476 \tex_mathchardef:D \c_max_register_int = 32767 ~
1477 \tex_fi:D
1478 \tex_fi:D

```

(End definition for `\c_max_register_int`. This variable is documented on page 100.)

5.3 Defining functions

We start by providing functions for the typical definition functions. First the local ones.

<pre> \cs_set_nopar:Npn \cs_set_nopar:Npx \cs_set:Npn \cs_set:Npx \cs_set_protected_nopar:Npn \cs_set_protected_nopar:Npx \cs_set_protected:Npn \cs_set_protected:Npx </pre>	<p>All assignment functions in L^AT_EX3 should be naturally protected; after all, the T_EX primitives for assignments are and it can be a cause of problems if others aren't.</p> <pre> 1479 \tex_let:D \cs_set_nopar:Npn \tex_def:D 1480 \tex_let:D \cs_set_nopar:Npx \tex_edef:D 1481 \tex_protected:D \tex_long:D \tex_def:D \cs_set:Npn 1482 { \tex_long:D \tex_def:D } 1483 \tex_protected:D \tex_long:D \tex_def:D \cs_set:Npx 1484 { \tex_long:D \tex_edef:D } 1485 \tex_protected:D \tex_long:D \tex_def:D \cs_set_protected_nopar:Npn 1486 { \tex_protected:D \tex_def:D } 1487 \tex_protected:D \tex_long:D \tex_def:D \cs_set_protected_nopar:Npx 1488 { \tex_protected:D \tex_edef:D } 1489 \tex_protected:D \tex_long:D \tex_def:D \cs_set_protected:Npn 1490 { \tex_protected:D \tex_long:D \tex_def:D } 1491 \tex_protected:D \tex_long:D \tex_def:D \cs_set_protected:Npx 1492 { \tex_protected:D \tex_long:D \tex_edef:D } </pre>
--	--

(End definition for `\cs_set_nopar:Npn` and others. These functions are documented on page 11.)

<pre> \cs_gset_nopar:Npn \cs_gset_nopar:Npx \cs_gset:Npn \cs_gset:Npx \cs_gset_protected_nopar:Npn \cs_gset_protected_nopar:Npx \cs_gset_protected:Npn \cs_gset_protected:Npx </pre>	<p>Global versions of the above functions.</p> <pre> 1493 \tex_let:D \cs_gset_nopar:Npn \tex_gdef:D 1494 \tex_let:D \cs_gset_nopar:Npx \tex_xdef:D 1495 \cs_set_protected:Npn \cs_gset:Npn 1496 { \tex_long:D \tex_gdef:D } 1497 \cs_set_protected:Npn \cs_gset:Npx 1498 { \tex_long:D \tex_xdef:D } 1499 \cs_set_protected:Npn \cs_gset_protected_nopar:Npn 1500 { \tex_protected:D \tex_gdef:D } 1501 \cs_set_protected:Npn \cs_gset_protected_nopar:Npx 1502 { \tex_protected:D \tex_xdef:D } 1503 \cs_set_protected:Npn \cs_gset_protected:Npn 1504 { \tex_protected:D \tex_long:D \tex_gdef:D } 1505 \cs_set_protected:Npn \cs_gset_protected:Npx 1506 { \tex_protected:D \tex_long:D \tex_xdef:D } </pre>
--	---

(End definition for `\cs_gset_nopar:Npn` and others. These functions are documented on page 12.)

5.4 Selecting tokens

<pre> 1507 <@@=exp> \l__exp_internal_tl </pre>	<p>Scratch token list variable for l3expan, used by <code>\use:x</code>, used in defining conditionals. We don't use <code>tl</code> methods because l3basics is loaded earlier.</p> <pre> 1508 \cs_set_nopar:Npn \l__exp_internal_tl { } </pre>
--	--

(End definition for `\l__exp_internal_tl`.)

`\use:c` This macro grabs its argument and returns a csname from it.

```
1509 \cs_set:Npn \use:c #1 { \cs:w #1 \cs_end: }
```

(End definition for `\use:c`. This function is documented on page 16.)

`\use:x` Fully expands its argument and passes it to the input stream. Uses the reserved `\l__exp_internal_tl` which we’ve set up above.

```
1510 \cs_set_protected:Npn \use:x #1
1511 {
1512   \cs_set_nopar:Npx \l__exp_internal_tl {#1}
1513   \l__exp_internal_tl
1514 }
```

(End definition for `\use:x`. This function is documented on page 20.)

```
1515 <@@=use>
```

`\use:e` In non-Lua \TeX engines older than 2019, `\expanded` is emulated.

```
1516 \cs_set:Npn \use:e #1 { \tex_expanded:D {#1} }
1517 \tex_ifdefined:D \tex_expanded:D \tex_else:D
1518   \cs_set:Npn \use:e #1 { \exp_args:Ne \use:n {#1} }
1519 \tex_fi:D
```

(End definition for `\use:e`. This function is documented on page 20.)

```
1520 <@@=exp>
```

`\use:n` These macros grab their arguments and return them back to the input (with outer braces removed).

```
\use:nnn 1521 \cs_set:Npn \use:n #1 {#1}
\use:nnnn 1522 \cs_set:Npn \use:nn #1#2 {#1#2}
1523 \cs_set:Npn \use:nnn #1#2#3 {#1#2#3}
1524 \cs_set:Npn \use:nnnn #1#2#3#4 {#1#2#3#4}
```

(End definition for `\use:n` and others. These functions are documented on page 19.)

`\use_i:nn` The equivalent to L^AT_EX 2_ε’s `\@firstoftwo` and `\@secondoftwo`.

```
\use_ii:nn 1525 \cs_set:Npn \use_i:nn #1#2 {#1}
1526 \cs_set:Npn \use_ii:nn #1#2 {#2}
```

(End definition for `\use_i:nn` and `\use_ii:nn`. These functions are documented on page 19.)

`\use_i:nnn` We also need something for picking up arguments from a longer list.

```
\use_ii:nnn 1527 \cs_set:Npn \use_i:nnn #1#2#3 {#1}
\use_iii:nnn 1528 \cs_set:Npn \use_ii:nnn #1#2#3 {#2}
\use_i_ii:nnn 1529 \cs_set:Npn \use_iii:nnn #1#2#3 {#3}
\use_i:nnnn 1530 \cs_set:Npn \use_i_ii:nnn #1#2#3 {#1#2}
\use_ii:nnnn 1531 \cs_set:Npn \use_i:nnnn #1#2#3#4 {#1}
\use_iii:nnnn 1532 \cs_set:Npn \use_ii:nnnn #1#2#3#4 {#2}
\use_iv:nnnn 1533 \cs_set:Npn \use_iii:nnnn #1#2#3#4 {#3}
1534 \cs_set:Npn \use_iv:nnnn #1#2#3#4 {#4}
```

(End definition for `\use_i:nnn` and others. These functions are documented on page 19.)

`\use_ii_i:nn`

```
1535 \cs_set:Npn \use_ii_i:nn #1#2 { #2 #1 }
```

(End definition for `\use_ii_i:nn`. This function is documented on page 20.)

`\use_none_delimit_by_q_nil:w`
`\use_none_delimit_by_q_stop:w`
`\use_none_delimit_by_q_recursion_stop:w`

Functions that gobble everything until they see either `\q_nil`, `\q_stop`, or `\q_recursion_stop`, respectively.

```
1536 \cs_set:Npn \use_none_delimit_by_q_nil:w #1 \q_nil { }
1537 \cs_set:Npn \use_none_delimit_by_q_stop:w #1 \q_stop { }
1538 \cs_set:Npn \use_none_delimit_by_q_recursion_stop:w #1 \q_recursion_stop { }
```

(End definition for `\use_none_delimit_by_q_nil:w`, `\use_none_delimit_by_q_stop:w`, and `\use_none_delimit_by_q_recursion_stop:w`. These functions are documented on page 21.)

`\use_i_delimit_by_q_nil:nw`
`\use_i_delimit_by_q_stop:nw`
`\use_i_delimit_by_q_recursion_stop:nw`

Same as above but execute first argument after gobbling. Very useful when you need to skip the rest of a mapping sequence but want an easy way to control what should be expanded next.

```
1539 \cs_set:Npn \use_i_delimit_by_q_nil:nw #1#2 \q_nil {#1}
1540 \cs_set:Npn \use_i_delimit_by_q_stop:nw #1#2 \q_stop {#1}
1541 \cs_set:Npn \use_i_delimit_by_q_recursion_stop:nw
1542 #1#2 \q_recursion_stop {#1}
```

(End definition for `\use_i_delimit_by_q_nil:nw`, `\use_i_delimit_by_q_stop:nw`, and `\use_i_delimit_by_q_recursion_stop:nw`. These functions are documented on page 21.)

5.5 Gobbling tokens from input

`\use_none:n`
`\use_none:nn`
`\use_none:nnn`
`\use_none:nnnn`
`\use_none:nnnnn`
`\use_none:nnnnnn`
`\use_none:nnnnnnn`
`\use_none:nnnnnnnn`

To gobble tokens from the input we use a standard naming convention: the number of tokens gobbled is given by the number of n's following the `:` in the name. Although we could define functions to remove ten arguments or more using separate calls of `\use_none:nnnnn`, this is very non-intuitive to the programmer who will assume that expanding such a function once takes care of gobbling all the tokens in one go.

```
1543 \cs_set:Npn \use_none:n #1 { }
1544 \cs_set:Npn \use_none:nn #1#2 { }
1545 \cs_set:Npn \use_none:nnn #1#2#3 { }
1546 \cs_set:Npn \use_none:nnnn #1#2#3#4 { }
1547 \cs_set:Npn \use_none:nnnnn #1#2#3#4#5 { }
1548 \cs_set:Npn \use_none:nnnnnn #1#2#3#4#5#6 { }
1549 \cs_set:Npn \use_none:nnnnnnn #1#2#3#4#5#6#7 { }
1550 \cs_set:Npn \use_none:nnnnnnnn #1#2#3#4#5#6#7#8 { }
1551 \cs_set:Npn \use_none:nnnnnnnnn #1#2#3#4#5#6#7#8#9 { }
```

(End definition for `\use_none:n` and others. These functions are documented on page 20.)

5.6 Debugging and patching later definitions

```
1552 <@@=debug>
```

`__kernel_if_debug:TF`

A more meaningful test of whether debugging is enabled than messing up with guards. We can also more easily change the logic in one place then. This is needed primarily for deprecations.

```
1553 \cs_set_protected:Npn \__kernel_if_debug:TF #1#2 {#2}
```

(End definition for `__kernel_if_debug:TF`.)

`\debug_on:n` Stubs.

```
\debug_off:n 1554 \cs_set_protected:Npn \debug_on:n #1
1555 {
1556   \__kernel_msg_error:nnx { kernel } { enable-debug }
1557   { \tl_to_str:n { \debug_on:n {#1} } }
1558 }
1559 \cs_set_protected:Npn \debug_off:n #1
1560 {
1561   \__kernel_msg_error:nnx { kernel } { enable-debug }
1562   { \tl_to_str:n { \debug_off:n {#1} } }
1563 }
```

(End definition for `\debug_on:n` and `\debug_off:n`. These functions are documented on page 24.)

`\debug_suspend:`

`\debug_resume:`

```
1564 \cs_set_protected:Npn \debug_suspend: { }
1565 \cs_set_protected:Npn \debug_resume: { }
```

(End definition for `\debug_suspend:` and `\debug_resume:`. These functions are documented on page 24.)

`__kernel_deprecation_code:nn`

`\g__debug_deprecation_on_tl`

`\g__debug_deprecation_off_tl`

Some commands were more recently deprecated and not yet removed; only make these into errors if the user requests it. This relies on two token lists, filled up in `l3deprecation`.

```
1566 \cs_set_nopar:Npn \g__debug_deprecation_on_tl { }
1567 \cs_set_nopar:Npn \g__debug_deprecation_off_tl { }
1568 \cs_set_protected:Npn \__kernel_deprecation_code:nn #1#2
1569 {
1570   \tl_gput_right:Nn \g__debug_deprecation_on_tl {#1}
1571   \tl_gput_right:Nn \g__debug_deprecation_off_tl {#2}
1572 }
```

(End definition for `__kernel_deprecation_code:nn`, `\g__debug_deprecation_on_tl`, and `\g__debug_deprecation_off_tl`.)

5.7 Conditional processing and definitions

```
1573 <@@=prg>
```

Underneath any predicate function (`_p`) or other conditional forms (TF, etc.) is a built-in logic saying that it after all of the testing and processing must return the *state* which leaves \TeX in. Therefore, a simple user interface could be something like

```
\if_meaning:w #1#2
  \prg_return_true:
\else:
  \if_meaning:w #1#3
    \prg_return_true:
  \else:
    \prg_return_false:
\fi:
\fi:
```

Usually, a \TeX programmer would have to insert a number of `\exp_after:wN`s to ensure the state value is returned at exactly the point where the last conditional is finished. However, that obscures the code and forces the \TeX programmer to prove that he/she knows the $2^n - 1$ table. We therefore provide the simpler interface.

`\prg_return_true:` The idea here is that `\exp:w` expands fully any `\else:` and `\fi:` that are waiting to be discarded, before reaching the `\exp_end:` which leaves an empty expansion. The code can then leave either the first or second argument in the input stream. This means that all of the branching code has to contain at least two tokens: see how the logical tests are actually implemented to see this.

```

1574 \cs_set:Npn \prg_return_true:
1575   { \exp_after:wN \use_i:nn \exp:w }
1576 \cs_set:Npn \prg_return_false:
1577   { \exp_after:wN \use_ii:nn \exp:w}

```

An extended state space could be implemented by including a more elaborate function in place of `\use_i:nn/\use_ii:nn`. Provided two arguments are absorbed then the code would work.

(End definition for \prg_return_true: and \prg_return_false:. These functions are documented on page 107.)

`\prg_use_none_delimit_by_q_recursion_stop:w` Private version of `\use_none_delimit_by_q_recursion_stop:w`.

```

1578 \cs_set:Npn \__prg_use_none_delimit_by_q_recursion_stop:w
1579   #1 \q__prg_recursion_stop { }

```

(End definition for __prg_use_none_delimit_by_q_recursion_stop:w.)

`\prg_set_conditional:Npnn` The user functions for the types using parameter text from the programmer. The various functions only differ by which function is used for the assignment. For those `Npnn` type functions, we must grab the parameter text, reading everything up to a left brace before continuing. Then split the base function into name and signature, and feed `{\<name>}` `{\<signature>}` `\<boolean>` `{\<set or new>}` `{\<maybe protected>}` `{\<parameters>}` `{TF,...}` `{\<code>}` to the auxiliary function responsible for defining all conditionals. Note that `e` stands for expandable and `p` for protected.

```

1580 \cs_set_protected:Npn \prg_set_conditional:Npnn
1581   { \__prg_generate_conditional_parm:NNNpnn \cs_set:Npn e }
1582 \cs_set_protected:Npn \prg_new_conditional:Npnn
1583   { \__prg_generate_conditional_parm:NNNpnn \cs_new:Npn e }
1584 \cs_set_protected:Npn \prg_set_protected_conditional:Npnn
1585   { \__prg_generate_conditional_parm:NNNpnn \cs_set_protected:Npn p }
1586 \cs_set_protected:Npn \prg_new_protected_conditional:Npnn
1587   { \__prg_generate_conditional_parm:NNNpnn \cs_new_protected:Npn p }
1588 \cs_set_protected:Npn \__prg_generate_conditional_parm:NNNpnn #1#2#3#4#
1589   {
1590     \use:x
1591     {
1592       \__prg_generate_conditional:nnNNNnnn
1593       \cs_split_function:N #3
1594     }
1595     #1 #2 {#4}
1596   }

```

(End definition for \prg_set_conditional:Npnn and others. These functions are documented on page 105.)

`\prg_set_conditional:Nnn` The user functions for the types automatically inserting the correct parameter text based on the signature. The various functions only differ by which function is used for the assignment. Split the base function into name and signature. The second auxiliary

`__prg_generate_conditional_count:NNNnn`
`__prg_generate_conditional_count:nnNNNnn`

generates the parameter text from the number of letters in the signature. Then feed $\langle name \rangle$ $\langle signature \rangle$ $\langle boolean \rangle$ $\langle set \text{ or } new \rangle$ $\langle maybe \text{ protected} \rangle$ $\langle parameters \rangle$ $\{TF, \dots\}$ $\langle code \rangle$ to the auxiliary function responsible for defining all conditionals. If the $\langle signature \rangle$ has more than 9 letters, the definition is aborted since T_EX macros have at most 9 arguments. The erroneous case where the function name contains no colon is captured later.

```

1597 \cs_set_protected:Npn \prg_set_conditional:Nnn
1598   { \prg_generate_conditional_count:NNNnn \cs_set:Npn e }
1599 \cs_set_protected:Npn \prg_new_conditional:Nnn
1600   { \prg_generate_conditional_count:NNNnn \cs_new:Npn e }
1601 \cs_set_protected:Npn \prg_set_protected_conditional:Nnn
1602   { \prg_generate_conditional_count:NNNnn \cs_set_protected:Npn p }
1603 \cs_set_protected:Npn \prg_new_protected_conditional:Nnn
1604   { \prg_generate_conditional_count:NNNnn \cs_new_protected:Npn p }
1605 \cs_set_protected:Npn \prg_generate_conditional_count:NNNnn #1#2#3
1606   {
1607     \use:x
1608     {
1609       \prg_generate_conditional_count:nnNNNnn
1610       \cs_split_function:N #3
1611     }
1612     #1 #2
1613   }
1614 \cs_set_protected:Npn \prg_generate_conditional_count:nnNNNnn #1#2#3#4#5
1615   {
1616     \kernel_cs_parm_from_arg_count:nnF
1617     { \prg_generate_conditional:nnNNNnnn {#1} {#2} #3 #4 #5 }
1618     { \tl_count:n {#2} }
1619     {
1620       \kernel_msg_error:nxxx { kernel } { bad-number-of-arguments }
1621       { \token_to_str:c { #1 : #2 } }
1622       { \tl_count:n {#2} }
1623       \use_none:nn
1624     }
1625   }

```

(End definition for $\backslash\text{prg_set_conditional:Nnn}$ and others. These functions are documented on page 105.)

$\backslash\text{prg_generate_conditional:nnNNNnnn}$
 $\backslash\text{prg_generate_conditional:NNnnnnNw}$
 $\backslash\text{prg_generate_conditional_test:w}$
 $\backslash\text{prg_generate_conditional_fast:nw}$

The workhorse here is going through a list of desired forms, *i.e.*, p, TF, T and F. The first three arguments come from splitting up the base form of the conditional, which gives the name, signature and a boolean to signal whether or not there was a colon in the name. In the absence of a colon, we throw an error and don't define any conditional. The fourth and fifth arguments build up the defining function. The sixth is the parameters to use (possibly empty), the seventh is the list of forms to define, the eighth is the replacement text which we will augment when defining the forms. The use of $\backslash\text{tl_to_str:n}$ makes the later loop more robust.

A large number of our low-level conditionals look like $\langle code \rangle \backslash\text{prg_return_true:} \backslash\text{else:} \backslash\text{prg_return_false:} \backslash\text{fi:}$ so we optimize this special case by calling $\backslash\text{prg_generate_conditional_fast:nw} \langle code \rangle$. This passes $\backslash\text{use_i:nn}$ instead of $\backslash\text{use_i_ii:nnn}$ to functions such as $\backslash\text{prg_generate_p_form:wNNnnnnN}$.

```

1626 \cs_set_protected:Npn \prg_generate_conditional:nnNNNnnn #1#2#3#4#5#6#7#8
1627   {

```

```

1628 \if_meaning:w \c_false_bool #3
1629 \__kernel_msg_error:nxx { kernel } { missing-colon }
1630 { \token_to_str:c {#1} }
1631 \exp_after:wN \use_none:nn
1632 \fi:
1633 \use:x
1634 {
1635 \exp_not:N \__prg_generate_conditional:NNnnnnNw
1636 \exp_not:n { #4 #5 {#1} {#2} {#6} }
1637 \__prg_generate_conditional_test:w
1638 #8 \s__prg_mark
1639 \__prg_generate_conditional_fast:nw
1640 \prg_return_true: \else: \prg_return_false: \fi: \s__prg_mark
1641 \use_none:n
1642 \exp_not:n { {#8} \use_i_ii:nnn }
1643 \tl_to_str:n {#7}
1644 \exp_not:n { , \q__prg_recursion_tail , \q__prg_recursion_stop }
1645 }
1646 }
1647 \cs_set:Npn \__prg_generate_conditional_test:w
1648 #1 \prg_return_true: \else: \prg_return_false: \fi: \s__prg_mark #2
1649 { #2 {#1} }
1650 \cs_set:Npn \__prg_generate_conditional_fast:nw #1#2 \exp_not:n #3
1651 { \exp_not:n { {#1} \use_i:nn } }

```

Looping through the list of desired forms. First are six arguments and seventh is the form. Use the form to call the correct type. If the form does not exist, the `\use:c` construction results in `\relax`, and the error message is displayed (unless the form is empty, to allow for {T, , F}), then `\use_none:nnnnnnnn` cleans up. Otherwise, the error message is removed by the variant form.

```

1652 \cs_set_protected:Npn \__prg_generate_conditional:NNnnnnNw #1#2#3#4#5#6#7#8 ,
1653 {
1654 \if_meaning:w \q__prg_recursion_tail #8
1655 \exp_after:wN \__prg_use_none_delimit_by_q_recursion_stop:w
1656 \fi:
1657 \use:c { __prg_generate_ #8 _form:wNNnnnnN }
1658 \tl_if_empty:nF {#8}
1659 {
1660 \__kernel_msg_error:nxxx
1661 { kernel } { conditional-form-unknown }
1662 {#8} { \token_to_str:c { #3 : #4 } }
1663 }
1664 \use_none:nnnnnnnn
1665 \s__prg_stop
1666 #1 #2 {#3} {#4} {#5} {#6} #7
1667 \__prg_generate_conditional:NNnnnnNw #1 #2 {#3} {#4} {#5} {#6} #7
1668 }

```

(End definition for `__prg_generate_conditional:nnNNnnnn` and others.)

```

\__prg_generate_p_form:wNNnnnnN
\__prg_generate_TF_form:wNNnnnnN
\__prg_generate_T_form:wNNnnnnN
\__prg_generate_F_form:wNNnnnnN
\__prg_p_true:w

```

How to generate the various forms. Those functions take the following arguments: 1: junk, 2: `\cs_set:Npn` or similar, 3: p (for protected conditionals) or e, 4: function name, 5: signature, 6: parameter text, 7: replacement (possibly trimmed by `__prg_generate_conditional_fast:nw`), 8: `\use_i_ii:nnn` or `\use_i:nn` (for “fast” conditionals). Remember that the logic-returning functions expect two arguments to be present

after `\exp_end::` notice the construction of the different variants relies on this, and that the `TF` and `F` variants will be slightly faster than the `T` version. The `p` form is only valid for expandable tests, we check for that by making sure that the second argument is empty. For “fast” conditionals, `#7` has an extra `\if_...` To optimize a bit further we could replace `\exp_after:wN \use_ii:nnn` and similar by a single macro similar to `__prg_p_true:w`. The drawback is that if the `T` or `F` arguments are actually missing, the recovery from the runaway argument would not insert `\fi:` back, messing up nesting of conditionals.

```

1669 \cs_set_protected:Npn \__prg_generate_p_form:wNNnnnnN
1670   #1 \s__prg_stop #2#3#4#5#6#7#8
1671   {
1672     \if_meaning:w e #3
1673     \exp_after:wN \use_i:nn
1674   \else:
1675     \exp_after:wN \use_ii:nn
1676   \fi:
1677     {
1678       #8
1679       { \exp_args:Nc #2 { #4 _p: #5 } #6 }
1680       { { #7 \exp_end: \c_true_bool \c_false_bool } }
1681       { #7 \__prg_p_true:w \fi: \c_false_bool }
1682     }
1683     {
1684       \__kernel_msg_error:nxx { kernel } { protected-predicate }
1685       { \token_to_str:c { #4 _p: #5 } }
1686     }
1687   }
1688 \cs_set_protected:Npn \__prg_generate_T_form:wNNnnnnN
1689   #1 \s__prg_stop #2#3#4#5#6#7#8
1690   {
1691     #8
1692     { \exp_args:Nc #2 { #4 : #5 T } #6 }
1693     { { #7 \exp_end: \use:n \use_none:n } }
1694     { #7 \exp_after:wN \use_ii:nn \fi: \use_none:n }
1695   }
1696 \cs_set_protected:Npn \__prg_generate_F_form:wNNnnnnN
1697   #1 \s__prg_stop #2#3#4#5#6#7#8
1698   {
1699     #8
1700     { \exp_args:Nc #2 { #4 : #5 F } #6 }
1701     { { #7 \exp_end: { } } }
1702     { #7 \exp_after:wN \use_none:nn \fi: \use:n }
1703   }
1704 \cs_set_protected:Npn \__prg_generate_TF_form:wNNnnnnN
1705   #1 \s__prg_stop #2#3#4#5#6#7#8
1706   {
1707     #8
1708     { \exp_args:Nc #2 { #4 : #5 TF } #6 }
1709     { { #7 \exp_end: } }
1710     { #7 \exp_after:wN \use_ii:nnn \fi: \use_ii:nn }
1711   }
1712 \cs_set:Npn \__prg_p_true:w \fi: \c_false_bool { \fi: \c_true_bool }

```

(End definition for `__prg_generate_p_form:wNNnnnnN` and others.)

`\prg_set_eq_conditional:NNn` The setting-equal functions. Split both functions and feed $\{\langle name_1 \rangle\}$ $\{\langle signature_1 \rangle\}$
`\prg_new_eq_conditional:NNn` $\langle boolean_1 \rangle$ $\{\langle name_2 \rangle\}$ $\{\langle signature_2 \rangle\}$ $\langle boolean_2 \rangle$ $\langle copying\ function \rangle$ $\langle conditions \rangle$, `\q__prg_set_eq_conditional:NNNn` `\prg_recursion_tail` , `\q__prg_recursion_stop` to a first auxiliary.

```

1713 \cs_set_protected:Npn \prg_set_eq_conditional:NNn
1714   { \__prg_set_eq_conditional:NNNn \cs_set_eq:cc }
1715 \cs_set_protected:Npn \prg_new_eq_conditional:NNn
1716   { \__prg_set_eq_conditional:NNNn \cs_new_eq:cc }
1717 \cs_set_protected:Npn \__prg_set_eq_conditional:NNNn #1#2#3#4
1718   {
1719     \use:x
1720     {
1721       \exp_not:N \__prg_set_eq_conditional:nnNnnNNw
1722       \cs_split_function:N #2
1723       \cs_split_function:N #3
1724       \exp_not:N #1
1725       \tl_to_str:n {#4}
1726       \exp_not:n { , \q__prg_recursion_tail , \q__prg_recursion_stop }
1727     }
1728   }

```

(End definition for `\prg_set_eq_conditional:NNn`, `\prg_new_eq_conditional:NNn`, and `__prg_set_eq_conditional:NNNn`. These functions are documented on page 106.)

`__prg_set_eq_conditional:nnNnnNNw`
`__prg_set_eq_conditional_loop:nnnnNw`
`__prg_set_eq_conditional_p_form:nnn`
`__prg_set_eq_conditional_TF_form:nnn`
`__prg_set_eq_conditional_T_form:nnn`
`__prg_set_eq_conditional_F_form:nnn`

Split the function to be defined, and setup a manual clist loop over argument #6 of the first auxiliary. The second auxiliary receives twice three arguments coming from splitting the function to be defined and the function to copy. Make sure that both functions contained a colon, otherwise we don't know how to build conditionals, hence abort. Call the looping macro, with arguments $\{\langle name_1 \rangle\}$ $\{\langle signature_1 \rangle\}$ $\{\langle name_2 \rangle\}$ $\{\langle signature_2 \rangle\}$ $\langle copying\ function \rangle$ and followed by the comma list. At each step in the loop, make sure that the conditional form we copy is defined, and copy it, otherwise abort.

```

1729 \cs_set_protected:Npn \__prg_set_eq_conditional:nnNnnNNw #1#2#3#4#5#6
1730   {
1731     \if_meaning:w \c_false_bool #3
1732       \__kernel_msg_error:nnx { kernel } { missing-colon }
1733       { \token_to_str:c {#1} } }
1734     \exp_after:wN \__prg_use_none_delimit_by_q_recursion_stop:w
1735     \fi:
1736     \if_meaning:w \c_false_bool #6
1737       \__kernel_msg_error:nnx { kernel } { missing-colon }
1738       { \token_to_str:c {#4} } }
1739     \exp_after:wN \__prg_use_none_delimit_by_q_recursion_stop:w
1740     \fi:
1741     \__prg_set_eq_conditional_loop:nnnnNw {#1} {#2} {#4} {#5}
1742   }
1743 \cs_set_protected:Npn \__prg_set_eq_conditional_loop:nnnnNw #1#2#3#4#5#6 ,
1744   {
1745     \if_meaning:w \q__prg_recursion_tail #6
1746       \exp_after:wN \__prg_use_none_delimit_by_q_recursion_stop:w
1747       \fi:
1748     \use:c { __prg_set_eq_conditional_ #6 _form:wNnnnn }
1749     \tl_if_empty:nF {#6}
1750     {
1751       \__kernel_msg_error:nnxx

```



```

1752         { kernel } { conditional-form-unknown }
1753         {#6} { \token_to_str:c { #1 : #2 } }
1754     }
1755     \use_none:nnnnnn
1756     \s__prg_stop
1757     #5 {#1} {#2} {#3} {#4}
1758     \__prg_set_eq_conditional_loop:nnnnNw {#1} {#2} {#3} {#4} #5
1759 }
1760 \cs_set:Npn \__prg_set_eq_conditional_p_form:wNnnnn #1 \s__prg_stop #2#3#4#5#6
1761 { #2 { #3 _p : #4 } { #5 _p : #6 } }
1762 \cs_set:Npn \__prg_set_eq_conditional_TF_form:wNnnnn #1 \s__prg_stop #2#3#4#5#6
1763 { #2 { #3 : #4 TF } { #5 : #6 TF } }
1764 \cs_set:Npn \__prg_set_eq_conditional_T_form:wNnnnn #1 \s__prg_stop #2#3#4#5#6
1765 { #2 { #3 : #4 T } { #5 : #6 T } }
1766 \cs_set:Npn \__prg_set_eq_conditional_F_form:wNnnnn #1 \s__prg_stop #2#3#4#5#6
1767 { #2 { #3 : #4 F } { #5 : #6 F } }

```

(End definition for `__prg_set_eq_conditional:nnNnnNNw` and others.)

All that is left is to define the canonical boolean true and false. I think Michael originated the idea of expandable boolean tests. At first these were supposed to expand into either TT or TF to be tested using `\if:w` but this was later changed to 00 and 01, so they could be used in logical operations. Later again they were changed to being numerical constants with values of 1 for true and 0 for false. We need this from the get-go.

`\c_true_bool` Here are the canonical boolean values.

`\c_false_bool`

```

1768 \tex_chardef:D \c_true_bool = 1 ~
1769 \tex_chardef:D \c_false_bool = 0 ~

```

(End definition for `\c_true_bool` and `\c_false_bool`. These variables are documented on page 22.)

5.8 Dissecting a control sequence

```

1770 <@@=cs>

```

`__cs_count_signature:N` `__cs_count_signature:N <function>`

Splits the `<function>` into the `<name>` (*i.e.* the part before the colon) and the `<signature>` (*i.e.* after the colon). The `<number>` of tokens in the `<signature>` is then left in the input stream. If there was no `<signature>` then the result is the marker value `-1`.

`__cs_get_function_name:N *` `__cs_get_function_name:N <function>`

Splits the `<function>` into the `<name>` (*i.e.* the part before the colon) and the `<signature>` (*i.e.* after the colon). The `<name>` is then left in the input stream without the escape character present made up of tokens with category code 12 (other).

`__cs_get_function_signature:N *` `__cs_get_function_signature:N <function>`

Splits the `<function>` into the `<name>` (*i.e.* the part before the colon) and the `<signature>` (*i.e.* after the colon). The `<signature>` is then left in the input stream made up of tokens with category code 12 (other).

<code>__cs_tmp:w</code>	Function used for various short-term usages, for instance defining functions whose definition involves tokens which are hard to insert normally (spaces, characters with category other).
<code>\cs_to_str:N</code> <code>__cs_to_str:N</code> <code>__cs_to_str:w</code>	This converts a control sequence into the character string of its name, removing the leading escape character. This turns out to be a non-trivial matter as there are different cases:

- The usual case of a printable escape character;
- the case of a non-printable escape characters, e.g., when the value of the `\escapechar` is negative;
- when the escape character is a space.

One approach to solve this is to test how many tokens result from `\token_to_str:N \a`. If there are two tokens, then the escape character is printable, while if it is non-printable then only one is present.

However, there is an additional complication: the control sequence itself may start with a space. Clearly that should *not* be lost in the process of converting to a string. So the approach adopted is a little more intricate still. When the escape character is printable, `\token_to_str:N _` yields the escape character itself and a space. The character codes are different, thus the `\if:w` test is false, and TeX reads `__cs_to_str:N` after turning the following control sequence into a string; this auxiliary removes the escape character, and stops the expansion of the initial `\tex_romannumeral:D`. The second case is that the escape character is not printable. Then the `\if:w` test is unfinished after reading a the space from `\token_to_str:N _`, and the auxiliary `__cs_to_str:w` is expanded, feeding – as a second character for the test; the test is false, and TeX skips to `\fi:`, then performs `\token_to_str:N`, and stops the `\tex_romannumeral:D` with `\c_zero_int`. The last case is that the escape character is itself a space. In this case, the `\if:w` test is true, and the auxiliary `__cs_to_str:w` comes into play, inserting `-\int_value:w`, which expands `\c_zero_int` to the character 0. The initial `\tex_romannumeral:D` then sees 0, which is not a terminated number, followed by the escape character, a space, which is removed, terminating the expansion of `\tex_romannumeral:D`. In all three cases, `\cs_to_str:N` takes two expansion steps to be fully expanded.

```

1771 \cs_set:Npn \cs_to_str:N
1772   {

```

We implement the expansion scheme using `\tex_romannumeral:D` terminating it with `\c_zero_int` rather than using `\exp:w` and `\exp_end:` as we normally do. The reason is that the code heavily depends on terminating the expansion with `\c_zero_int` so we make this dependency explicit.

```

1773   \tex_romannumeral:D
1774   \if:w \token_to_str:N \ \__cs_to_str:w \fi:
1775   \exp_after:wN \__cs_to_str:N \token_to_str:N
1776 }
1777 \cs_set:Npn \__cs_to_str:N #1 { \c_zero_int }
1778 \cs_set:Npn \__cs_to_str:w #1 \__cs_to_str:N
1779 { - \int_value:w \fi: \exp_after:wN \c_zero_int }

```

If speed is a concern we could use `\csstring` in LuaTeX. For the empty csname that primitive gives an empty result while the current `\cs_to_str:N` gives incorrect results in all engines (this is impossible to fix without huge performance hit).

(End definition for `\cs_to_str:N`, `__cs_to_str:N`, and `__cs_to_str:w`. This function is documented on page 17.)

`\cs_split_function:N`

This function takes a function name and splits it into name with the escape char removed and argument specification. In addition to this, a third argument, a boolean $\langle true \rangle$ or $\langle false \rangle$ is returned with $\langle true \rangle$ for when there is a colon in the function and $\langle false \rangle$ if there is not.

We cannot use `:` directly as it has the wrong category code so an `x`-type expansion is used to force the conversion.

First ensure that we actually get a properly evaluated string by expanding `\cs_to_str:N` twice. If the function contained a colon, the auxiliary takes as `#1` the function name, delimited by the first colon, then the signature `#2`, delimited by `\s__cs_mark`, then `\c_true_bool` as `#3`, and `#4` cleans up until `\s__cs_stop`. Otherwise, the `#1` contains the function name and `\s__cs_mark` `\c_true_bool`, `#2` is empty, `#3` is `\c_false_bool`, and `#4` cleans up. The second auxiliary trims the trailing `\s__cs_mark` from the function name if present (that is, if the original function had no colon).

```

1780 \cs_set_protected:Npn \__cs_tmp:w #1
1781 {
1782   \cs_set:Npn \cs_split_function:N ##1
1783   {
1784     \exp_after:wN \exp_after:wN \exp_after:wN
1785     \__cs_split_function_auxi:w
1786     \cs_to_str:N ##1 \s__cs_mark \c_true_bool
1787     #1 \s__cs_mark \c_false_bool \s__cs_stop
1788   }
1789   \cs_set:Npn \__cs_split_function_auxi:w
1790     ##1 #1 ##2 \s__cs_mark ##3##4 \s__cs_stop
1791     { \__cs_split_function_auxii:w ##1 \s__cs_mark \s__cs_stop {##2} ##3 }
1792   \cs_set:Npn \__cs_split_function_auxii:w ##1 \s__cs_mark ##2 \s__cs_stop
1793     { {##1} }
1794 }
1795 \exp_after:wN \__cs_tmp:w \token_to_str:N :
```

(End definition for `\cs_split_function:N`, `__cs_split_function_auxi:w`, and `__cs_split_function_auxii:w`. This function is documented on page 17.)

5.9 Exist or free

A control sequence is said to *exist* (to be used) if has an entry in the hash table and its meaning is different from the primitive `\relax` token. A control sequence is said to be *free* (to be defined) if it does not already exist.

`\cs_if_exist_p:N`

`\cs_if_exist_p:c`

`\cs_if_exist:NTF`

`\cs_if_exist:cTF`

Two versions for checking existence. For the `N` form we firstly check for `\scan_stop:` and then if it is in the hash table. There is no problem when inputting something like `\else:` or `\fi:` as \TeX will only ever skip input in case the token tested against is `\scan_stop:`.

```

1796 \prg_set_conditional:Npnn \cs_if_exist:N #1 { p , T , F , TF }
1797 {
1798   \if_meaning:w #1 \scan_stop:
1799   \prg_return_false:
1800   \else:
1801     \if_cs_exist:N #1
1802     \prg_return_true:
1803     \else:
```

```

1804     \prg_return_false:
1805     \fi:
1806   \fi:
1807 }

```

For the `c` form we firstly check if it is in the hash table and then for `\scan_stop`: so that we do not add it to the hash table unless it was already there. Here we have to be careful as the text to be skipped if the first test is false may contain tokens that disturb the scanner. Therefore, we ensure that the second test is performed after the first one has concluded completely.

```

1808 \prg_set_conditional:Npnn \cs_if_exist:c #1 { p , T , F , TF }
1809 {
1810   \if_cs_exist:w #1 \cs_end:
1811   \exp_after:wN \use_i:nn
1812   \else:
1813     \exp_after:wN \use_ii:nn
1814   \fi:
1815   {
1816     \exp_after:wN \if_meaning:w \cs:w #1 \cs_end: \scan_stop:
1817     \prg_return_false:
1818     \else:
1819       \prg_return_true:
1820     \fi:
1821   }
1822   \prg_return_false:
1823 }

```

(End definition for `\cs_if_exist:NTF`. This function is documented on page 22.)

`\cs_if_free_p:N`
`\cs_if_free_p:c`
`\cs_if_free:NTF`
`\cs_if_free:cTF`

The logical reversal of the above.

```

1824 \prg_set_conditional:Npnn \cs_if_free:N #1 { p , T , F , TF }
1825 {
1826   \if_meaning:w #1 \scan_stop:
1827   \prg_return_true:
1828   \else:
1829     \if_cs_exist:N #1
1830     \prg_return_false:
1831     \else:
1832       \prg_return_true:
1833     \fi:
1834   \fi:
1835 }
1836 \prg_set_conditional:Npnn \cs_if_free:c #1 { p , T , F , TF }
1837 {
1838   \if_cs_exist:w #1 \cs_end:
1839   \exp_after:wN \use_i:nn
1840   \else:
1841     \exp_after:wN \use_ii:nn
1842   \fi:
1843   {
1844     \exp_after:wN \if_meaning:w \cs:w #1 \cs_end: \scan_stop:
1845     \prg_return_true:
1846     \else:
1847       \prg_return_false:

```

```

1848     \fi:
1849   }
1850   { \prg_return_true: }
1851 }

```

(End definition for `\cs_if_free:NTF`. This function is documented on page 22.)

`\cs_if_exist_use:N` The `\cs_if_exist_use:...` functions cannot be implemented as conditionals because the true branch must leave both the control sequence itself and the true code in the input stream. For the `c` variants, we are careful not to put the control sequence in the hash table if it does not exist. In LuaTeX we could use the `\lastnamedcs` primitive.

```

1852 \cs_set:Npn \cs_if_exist_use:NTF #1#2
1853 { \cs_if_exist:NTF #1 { #1 #2 } }
1854 \cs_set:Npn \cs_if_exist_use:NF #1
1855 { \cs_if_exist:NTF #1 { #1 } }
1856 \cs_set:Npn \cs_if_exist_use:NT #1 #2
1857 { \cs_if_exist:NTF #1 { #1 #2 } { } }
1858 \cs_set:Npn \cs_if_exist_use:N #1
1859 { \cs_if_exist:NTF #1 { #1 } { } }
1860 \cs_set:Npn \cs_if_exist_use:cTF #1#2
1861 { \cs_if_exist:cTF {#1} { \use:c {#1} #2 } }
1862 \cs_set:Npn \cs_if_exist_use:cF #1
1863 { \cs_if_exist:cTF {#1} { \use:c {#1} } }
1864 \cs_set:Npn \cs_if_exist_use:cT #1#2
1865 { \cs_if_exist:cTF {#1} { \use:c {#1} #2 } { } }
1866 \cs_set:Npn \cs_if_exist_use:c #1
1867 { \cs_if_exist:cTF {#1} { \use:c {#1} } { } }

```

(End definition for `\cs_if_exist_use:NTF`. This function is documented on page 16.)

5.10 Preliminaries for new functions

We provide two kinds of functions that can be used to define control sequences. On the one hand we have functions that check if their argument doesn't already exist, they are called `\..._new`. The second type of defining functions doesn't check if the argument is already defined.

Before we can define them, we need some auxiliary macros that allow us to generate error messages. The next few definitions here are only temporary, they will be redefined later on.

`__kernel_msg_error:nxxx` If an internal error occurs before L^AT_EX3 has loaded `l3msg` then the code should issue a usable if terse error message and halt. This can only happen if a coding error is made by the team, so this is a reasonable response. Setting the `\newlinechar` is needed, to turn `^^J` into a proper line break in plain T_EX.

```

1868 \cs_set_protected:Npn \__kernel_msg_error:nxxx #1#2#3#4
1869 {
1870   \tex_newlinechar:D = '\^^J \scan_stop:
1871   \tex_errmessage:D
1872   {
1873     !!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!!~! ^^J
1874     Argh,~internal~LaTeX3~error! ^^J ^^J
1875     Module ~ #1 , ~ message~name~"#2": ^^J
1876     Arguments~'#3'~and~'#4' ^^J ^^J

```

```

1877         This~is~one~for~The~LaTeX3~Project:~bailing~out
1878     }
1879     \tex_end:D
1880 }
1881 \cs_set_protected:Npn \__kernel_msg_error:nxx #1#2#3
1882 { \__kernel_msg_error:nxxx {#1} {#2} {#3} { } }
1883 \cs_set_protected:Npn \__kernel_msg_error:nn #1#2
1884 { \__kernel_msg_error:nxxx {#1} {#2} { } { } }

```

(End definition for `__kernel_msg_error:nxxx`, `__kernel_msg_error:nxx`, and `__kernel_msg_error:nn`.)

`\msg_line_context:` Another one from `l3msg` which will be altered later.

```

1885 \cs_set:Npn \msg_line_context:
1886 { on~line~ \tex_the:D \tex_inputlineno:D }

```

(End definition for `\msg_line_context:`. This function is documented on page 154.)

`\iow_log:x` We define a routine to write only to the log file. And a similar one for writing to both
`\iow_term:x` the log file and the terminal. These will be redefined later by `l3io`.

```

1887 \cs_set_protected:Npn \iow_log:x
1888 { \tex_immediate:D \tex_write:D -1 }
1889 \cs_set_protected:Npn \iow_term:x
1890 { \tex_immediate:D \tex_write:D 16 }

```

(End definition for `\iow_log:n`. This function is documented on page 164.)

`__kernel_chk_if_free_cs:N` This command is called by `\cs_new_nopar:Npn` and `\cs_new_eq:NN` etc. to make sure
`__kernel_chk_if_free_cs:c` that the argument sequence is not already in use. If it is, an error is signalled. It checks
if `<csname>` is undefined or `\scan_stop:`. Otherwise an error message is issued. We have
to make sure we don't put the argument into the conditional processing since it may be
an `\if...` type function!

```

1891 \cs_set_protected:Npn \__kernel_chk_if_free_cs:N #1
1892 {
1893     \cs_if_free:NF #1
1894     {
1895         \__kernel_msg_error:nxxx { kernel } { command-already-defined }
1896         { \token_to_str:N #1 } { \token_to_meaning:N #1 }
1897     }
1898 }
1899 \cs_set_protected:Npn \__kernel_chk_if_free_cs:c
1900 { \exp_args:Nc \__kernel_chk_if_free_cs:N }

```

(End definition for `__kernel_chk_if_free_cs:N`.)

5.11 Defining new functions

```

1901 (@@=cs)

```

`\cs_new_nopar:Npn` Function which check that the control sequence is free before defining it.

```

\cs_new_nopar:Npx
\cs_new:Npn
\cs_new:Npx
\cs_new_protected_nopar:Npn
\cs_new_protected_nopar:Npx
\cs_new_protected:Npn
\cs_new_protected:Npx
\__cs_tmp:w

```

```

1902 \cs_set:Npn \__cs_tmp:w #1#2
1903 {
1904     \cs_set_protected:Npn #1 ##1
1905     {
1906         \__kernel_chk_if_free_cs:N ##1
1907         #2 ##1

```

```

1908     }
1909 }
1910 \__cs_tmp:w \cs_new_nopar:Npn          \cs_gset_nopar:Npn
1911 \__cs_tmp:w \cs_new_nopar:Npx          \cs_gset_nopar:Npx
1912 \__cs_tmp:w \cs_new:Npn                \cs_gset:Npn
1913 \__cs_tmp:w \cs_new:Npx                \cs_gset:Npx
1914 \__cs_tmp:w \cs_new_protected_nopar:Npn \cs_gset_protected_nopar:Npn
1915 \__cs_tmp:w \cs_new_protected_nopar:Npx \cs_gset_protected_nopar:Npx
1916 \__cs_tmp:w \cs_new_protected:Npn      \cs_gset_protected:Npn
1917 \__cs_tmp:w \cs_new_protected:Npx      \cs_gset_protected:Npx

```

(End definition for `\cs_new_nopar:Npn` and others. These functions are documented on page 11.)

`\cs_set_nopar:cpn` Like `\cs_set_nopar:Npn` and `\cs_new_nopar:Npn`, except that the first argument consists of the sequence of characters that should be used to form the name of the desired control sequence (the `c` stands for `csname` argument, see the expansion module). Global versions are also provided.

`\cs_set_nopar:cpn` $\langle string \rangle \langle rep-text \rangle$ turns $\langle string \rangle$ into a `csname` and then assigns $\langle rep-text \rangle$ to it by using `\cs_set_nopar:Npn`. This means that there might be a parameter string between the two arguments.

```

1918 \cs_set:Npn \__cs_tmp:w #1#2
1919 { \cs_new_protected_nopar:Npn #1 { \exp_args:Nc #2 } }
1920 \__cs_tmp:w \cs_set_nopar:cpn \cs_set_nopar:Npn
1921 \__cs_tmp:w \cs_set_nopar:cpx \cs_set_nopar:Npx
1922 \__cs_tmp:w \cs_gset_nopar:cpn \cs_gset_nopar:Npn
1923 \__cs_tmp:w \cs_gset_nopar:cpx \cs_gset_nopar:Npx
1924 \__cs_tmp:w \cs_new_nopar:cpn \cs_new_nopar:Npn
1925 \__cs_tmp:w \cs_new_nopar:cpx \cs_new_nopar:Npx

```

(End definition for `\cs_set_nopar:Npn`. This function is documented on page 11.)

`\cs_set:cpn` Variants of the `\cs_set:Npn` versions which make a `csname` out of the first arguments.
`\cs_set:cpx` We may also do this globally.

```

1926 \__cs_tmp:w \cs_set:cpn \cs_set:Npn
1927 \__cs_tmp:w \cs_set:cpx \cs_set:Npx
1928 \__cs_tmp:w \cs_gset:cpn \cs_gset:Npn
1929 \__cs_tmp:w \cs_gset:cpx \cs_gset:Npx
1930 \__cs_tmp:w \cs_new:cpn \cs_new:Npn
1931 \__cs_tmp:w \cs_new:cpx \cs_new:Npx

```

(End definition for `\cs_set:Npn`. This function is documented on page 11.)

`\cs_set_protected_nopar:cpn` Variants of the `\cs_set_protected_nopar:Npn` versions which make a `csname` out of the first arguments. We may also do this globally.

```

1932 \__cs_tmp:w \cs_set_protected_nopar:cpn \cs_set_protected_nopar:Npn
1933 \__cs_tmp:w \cs_set_protected_nopar:cpx \cs_set_protected_nopar:Npx
1934 \__cs_tmp:w \cs_gset_protected_nopar:cpn \cs_gset_protected_nopar:Npn
1935 \__cs_tmp:w \cs_gset_protected_nopar:cpx \cs_gset_protected_nopar:Npx
1936 \__cs_tmp:w \cs_new_protected_nopar:cpn \cs_new_protected_nopar:Npn
1937 \__cs_tmp:w \cs_new_protected_nopar:cpx \cs_new_protected_nopar:Npx

```

(End definition for `\cs_set_protected_nopar:Npn`. This function is documented on page 12.)

<code>\cs_set_protected:cpn</code>	1938	<code>__cs_tmp:w \cs_set_protected:cpn \cs_set_protected:Npn</code>
<code>\cs_set_protected:cpx</code>	1939	<code>__cs_tmp:w \cs_set_protected:cpx \cs_set_protected:Npx</code>
<code>\cs_gset_protected:cpn</code>	1940	<code>__cs_tmp:w \cs_gset_protected:cpn \cs_gset_protected:Npn</code>
<code>\cs_gset_protected:cpx</code>	1941	<code>__cs_tmp:w \cs_gset_protected:cpx \cs_gset_protected:Npx</code>
<code>\cs_new_protected:cpn</code>	1942	<code>__cs_tmp:w \cs_new_protected:cpn \cs_new_protected:Npn</code>
<code>\cs_new_protected:cpx</code>	1943	<code>__cs_tmp:w \cs_new_protected:cpx \cs_new_protected:Npx</code>

(End definition for `\cs_set_protected:Npn`. This function is documented on page 11.)

5.12 Copying definitions

<code>\cs_set_eq:NN</code>	These macros allow us to copy the definition of a control sequence to another control sequence.
<code>\cs_set_eq:cN</code>	The = sign allows us to define funny char tokens like = itself or <code>_</code> with this function.
<code>\cs_set_eq:Nc</code>	For the definition of <code>\c_space_char{~}</code> to work we need the ~ after the =.
<code>\cs_set_eq:cc</code>	
<code>\cs_gset_eq:NN</code>	<code>\cs_set_eq:NN</code> is long to avoid problems with a literal argument of <code>\par</code> . While
<code>\cs_gset_eq:cN</code>	<code>\cs_new_eq:NN</code> will probably never be correct with a first argument of <code>\par</code> , define it
<code>\cs_gset_eq:Nc</code>	long in order to throw an “already defined” error rather than “runaway argument”.
<code>\cs_gset_eq:cc</code>	1944 <code>\cs_new_protected:Npn \cs_set_eq:NN #1 { \tex_let:D #1 =~ }</code>
<code>\cs_new_eq:NN</code>	1945 <code>\cs_new_protected:Npn \cs_set_eq:cN { \exp_args:Nc \cs_set_eq:NN }</code>
<code>\cs_new_eq:cN</code>	1946 <code>\cs_new_protected:Npn \cs_set_eq:Nc { \exp_args:NNc \cs_set_eq:NN }</code>
<code>\cs_new_eq:Nc</code>	1947 <code>\cs_new_protected:Npn \cs_set_eq:cc { \exp_args:Ncc \cs_set_eq:NN }</code>
<code>\cs_new_eq:cc</code>	1948 <code>\cs_new_protected:Npn \cs_gset_eq:NN { \tex_global:D \cs_set_eq:NN }</code>
	1949 <code>\cs_new_protected:Npn \cs_gset_eq:Nc { \exp_args:NNc \cs_gset_eq:NN }</code>
	1950 <code>\cs_new_protected:Npn \cs_gset_eq:cN { \exp_args:Nc \cs_gset_eq:NN }</code>
	1951 <code>\cs_new_protected:Npn \cs_gset_eq:cc { \exp_args:Ncc \cs_gset_eq:NN }</code>
	1952 <code>\cs_new_protected:Npn \cs_new_eq:NN #1</code>
	1953 <code>{</code>
	1954 <code> __kernel_chk_if_free_cs:N #1</code>
	1955 <code> \tex_global:D \cs_set_eq:NN #1</code>
	1956 <code>}</code>
	1957 <code>\cs_new_protected:Npn \cs_new_eq:cN { \exp_args:Nc \cs_new_eq:NN }</code>
	1958 <code>\cs_new_protected:Npn \cs_new_eq:Nc { \exp_args:NNc \cs_new_eq:NN }</code>
	1959 <code>\cs_new_protected:Npn \cs_new_eq:cc { \exp_args:Ncc \cs_new_eq:NN }</code>

(End definition for `\cs_set_eq:NN`, `\cs_gset_eq:NN`, and `\cs_new_eq:NN`. These functions are documented on page 15.)

5.13 Undefined functions

<code>\cs_undefine:N</code>	The following function is used to free the main memory from the definition of some
<code>\cs_undefine:c</code>	function that isn't in use any longer. The c variant is careful not to add the control
	sequence to the hash table if it isn't there yet, and it also avoids nesting TeX conditionals
	in case #1 is unbalanced in this matter.

```

1960 \cs_new_protected:Npn \cs_undefine:N #1
1961 { \cs_gset_eq:NN #1 \tex_undefined:D }
1962 \cs_new_protected:Npn \cs_undefine:c #1
1963 {
1964   \if_cs_exist:w #1 \cs_end:
1965     \exp_after:wN \use:n
1966   \else:

```



```

1967     \exp_after:wN \use_none:n
1968   \fi:
1969   { \cs_gset_eq:cN {#1} \tex_undefined:D }
1970 }

```

(End definition for `\cs_undefine:N`. This function is documented on page 15.)

5.14 Generating parameter text from argument count

```

1971 <@@=cs>

```

```

\__kernel_cs_parm_from_arg_count:nnF
\__cs_parm_from_arg_count_test:nnF

```

LaTeX3 provides shorthands to define control sequences and conditionals with a simple parameter text, derived directly from the signature, or more generally from knowing the number of arguments, between 0 and 9. This function expands to its first argument, untouched, followed by a brace group containing the parameter text `{#1...#n}`, where n is the result of evaluating the second argument (as described in `\int_eval:n`). If the second argument gives a result outside the range $[0, 9]$, the third argument is returned instead, normally an error message. Some of the functions use here are not defined yet, but will be defined before this function is called.

```

1972 \cs_set_protected:Npn \__kernel_cs_parm_from_arg_count:nnF #1#2
1973 {
1974   \exp_args:Nx \__cs_parm_from_arg_count_test:nnF
1975   {
1976     \exp_after:wN \exp_not:n
1977     \if_case:w \int_eval:n {#2}
1978       { }
1979       \or: { ##1 }
1980       \or: { ##1##2 }
1981       \or: { ##1##2##3 }
1982       \or: { ##1##2##3##4 }
1983       \or: { ##1##2##3##4##5 }
1984       \or: { ##1##2##3##4##5##6 }
1985       \or: { ##1##2##3##4##5##6##7 }
1986       \or: { ##1##2##3##4##5##6##7##8 }
1987       \or: { ##1##2##3##4##5##6##7##8##9 }
1988       \else: { \c_false_bool }
1989     \fi:
1990   }
1991   {#1}
1992 }
1993 \cs_set_protected:Npn \__cs_parm_from_arg_count_test:nnF #1#2
1994 {
1995   \if_meaning:w \c_false_bool #1
1996   \exp_after:wN \use_ii:nn
1997   \else:
1998   \exp_after:wN \use_i:nn
1999   \fi:
2000   { #2 {#1} }
2001 }

```

(End definition for `__kernel_cs_parm_from_arg_count:nnF` and `__cs_parm_from_arg_count_test:nnF`.)

5.15 Defining functions from a given number of arguments

2002 `<@@=cs>`

`_cs_count_signature:N` Counting the number of tokens in the signature, *i.e.*, the number of arguments the function should take. Since this is not used in any time-critical function, we simply use `\tl_count:n` if there is a signature, otherwise `-1` arguments to signal an error. We need a variant form right away.

```

2003 \cs_new:Npn \_cs_count_signature:N #1
2004 { \exp_args:Nf \_cs_count_signature:n { \cs_split_function:N #1 } }
2005 \cs_new:Npn \_cs_count_signature:n #1
2006 { \int_eval:n { \_cs_count_signature:nnN #1 } }
2007 \cs_new:Npn \_cs_count_signature:nnN #1#2#3
2008 {
2009   \if_meaning:w \c_true_bool #3
2010     \tl_count:n {#2}
2011   \else:
2012     -1
2013   \fi:
2014 }
2015 \cs_new:Npn \_cs_count_signature:c
2016 { \exp_args:Nc \_cs_count_signature:N }
```

(End definition for `_cs_count_signature:N`, `_cs_count_signature:n`, and `_cs_count_signature:nnN`.)

`\cs_generate_from_arg_count:NNnn`
`\cs_generate_from_arg_count:cNnn`
`\cs_generate_from_arg_count:Ncnn`

We provide a constructor function for defining functions with a given number of arguments. For this we need to choose the correct parameter text and then use that when defining. Since TeX supports from zero to nine arguments, we use a simple switch to choose the correct parameter text, ensuring the result is returned after finishing the conditional. If it is not between zero and nine, we throw an error.

1: function to define, 2: with what to define it, 3: the number of args it requires and 4: the replacement text

```

2017 \cs_new_protected:Npn \cs_generate_from_arg_count:NNnn #1#2#3#4
2018 {
2019   \_kernel_cs_parm_from_arg_count:nnF { \use:nnn #2 #1 } {#3}
2020   {
2021     \_kernel_msg_error:nnxx { kernel } { bad-number-of-arguments }
2022     { \token_to_str:N #1 } { \int_eval:n {#3} }
2023     \use_none:n
2024   }
2025   {#4}
2026 }
```

A variant form we need right away, plus one which is used elsewhere but which is most logically created here.

```

2027 \cs_new_protected:Npn \cs_generate_from_arg_count:cNnn
2028 { \exp_args:Nc \cs_generate_from_arg_count:NNnn }
2029 \cs_new_protected:Npn \cs_generate_from_arg_count:Ncnn
2030 { \exp_args:NNc \cs_generate_from_arg_count:NNnn }
```

(End definition for `\cs_generate_from_arg_count:NNnn`. This function is documented on page 14.)

5.16 Using the signature to define functions

2031 <@@=cs>

We can now combine some of the tools we have to provide a simple interface for defining functions, where the number of arguments is read from the signature. For instance, `\cs_set:Nn \foo_bar:nn {#1,#2}`.

We want to define `\cs_set:Nn` as

```
\cs_set:Nn
\cs_set:Nx
\cs_set_nopar:Nn
\cs_set_nopar:Nx
\cs_set_protected:Nn
\cs_set_protected:Nx
\cs_set_protected_nopar:Nn
\cs_set_protected_nopar:Nx
\cs_gset:Nn
\cs_gset:Nx
\cs_gset_nopar:Nn
\cs_gset_nopar:Nx
\cs_gset_protected:Nn
\cs_gset_protected:Nx
\cs_gset_protected_nopar:Nn
\cs_gset_protected_nopar:Nx
\cs_new:Nn
\cs_new:Nx
\cs_new_nopar:Nn
\cs_new_nopar:Nx
\cs_new_protected:Nn
\cs_new_protected:Nx
\cs_new_protected_nopar:Nn
\cs_new_protected_nopar:Nx
```

```
\cs_set_protected:Npn \cs_set:Nn #1#2
{
  \cs_generate_from_arg_count:NNnn #1 \cs_set:Npn
  { \@@_count_signature:N #1 } {#2}
}
```

In short, to define `\cs_set:Nn` we need just use `\cs_set:Npn`, everything else is the same for each variant. Therefore, we can make it simpler by temporarily defining a function to do this for us.

```
2032 \cs_set:Npn \__cs_tmp:w #1#2#3
2033 {
2034   \cs_new_protected:cpx { cs_ #1 : #2 }
2035   {
2036     \exp_not:N \__cs_generate_from_signature:NNnn
2037     \exp_after:wN \exp_not:N \cs:w cs_ #1 : #3 \cs_end:
2038   }
2039 }
2040 \cs_new_protected:Npn \__cs_generate_from_signature:NNnn #1#2
2041 {
2042   \use:x
2043   {
2044     \__cs_generate_from_signature:nnNNnn
2045     \cs_split_function:N #2
2046   }
2047   #1 #2
2048 }
2049 \cs_new_protected:Npn \__cs_generate_from_signature:nnNNnn #1#2#3#4#5#6
2050 {
2051   \bool_if:NTF #3
2052   {
2053     \str_if_eq:eeF { }
2054     { \tl_map_function:nN {#2} \__cs_generate_from_signature:n }
2055     {
2056       \__kernel_msg_error:nnx { kernel } { non-base-function }
2057       { \token_to_str:N #5 }
2058     }
2059     \cs_generate_from_arg_count:NNnn
2060     #5 #4 { \tl_count:n {#2} } {#6}
2061   }
2062   {
2063     \__kernel_msg_error:nnx { kernel } { missing-colon }
2064     { \token_to_str:N #5 }
2065   }
2066 }
2067 \cs_new:Npn \__cs_generate_from_signature:n #1
```

```

2068 {
2069     \if:w n #1 \else: \if:w N #1 \else:
2070     \if:w T #1 \else: \if:w F #1 \else: #1 \fi: \fi: \fi: \fi:
2071 }

```

Then we define the 24 variants beginning with N.

```

2072 \__cs_tmp:w { set } { Nn } { Npn }
2073 \__cs_tmp:w { set } { Nx } { Npx }
2074 \__cs_tmp:w { set_nopar } { Nn } { Npn }
2075 \__cs_tmp:w { set_nopar } { Nx } { Npx }
2076 \__cs_tmp:w { set_protected } { Nn } { Npn }
2077 \__cs_tmp:w { set_protected } { Nx } { Npx }
2078 \__cs_tmp:w { set_protected_nopar } { Nn } { Npn }
2079 \__cs_tmp:w { set_protected_nopar } { Nx } { Npx }
2080 \__cs_tmp:w { gset } { Nn } { Npn }
2081 \__cs_tmp:w { gset } { Nx } { Npx }
2082 \__cs_tmp:w { gset_nopar } { Nn } { Npn }
2083 \__cs_tmp:w { gset_nopar } { Nx } { Npx }
2084 \__cs_tmp:w { gset_protected } { Nn } { Npn }
2085 \__cs_tmp:w { gset_protected } { Nx } { Npx }
2086 \__cs_tmp:w { gset_protected_nopar } { Nn } { Npn }
2087 \__cs_tmp:w { gset_protected_nopar } { Nx } { Npx }
2088 \__cs_tmp:w { new } { Nn } { Npn }
2089 \__cs_tmp:w { new } { Nx } { Npx }
2090 \__cs_tmp:w { new_nopar } { Nn } { Npn }
2091 \__cs_tmp:w { new_nopar } { Nx } { Npx }
2092 \__cs_tmp:w { new_protected } { Nn } { Npn }
2093 \__cs_tmp:w { new_protected } { Nx } { Npx }
2094 \__cs_tmp:w { new_protected_nopar } { Nn } { Npn }
2095 \__cs_tmp:w { new_protected_nopar } { Nx } { Npx }

```

(End definition for \cs_set:Nn and others. These functions are documented on page 13.)

\cs_set:cn The 24 c variants simply use \exp_args:Nc.

\cs_set:cx 2096 \cs_set:Npn __cs_tmp:w #1#2

\cs_set_nopar:cn 2097 {

\cs_set_nopar:cx 2098 \cs_new_protected:cpx { cs_ #1 : c #2 }

\cs_set_protected:cn 2099 {

\cs_set_protected:cx 2100 \exp_not:N \exp_args:Nc

\cs_set_protected_nopar:cn 2101 \exp_after:wN \exp_not:N \cs:w cs_ #1 : N #2 \cs_end:

\cs_set_protected_nopar:cx 2102 }

\cs_gset:cn 2103 }

\cs_gset:cx 2104 __cs_tmp:w { set } { n }

\cs_gset_nopar:cn 2105 __cs_tmp:w { set } { x }

\cs_gset_nopar:cx 2106 __cs_tmp:w { set_nopar } { n }

\cs_gset_protected:cn 2107 __cs_tmp:w { set_nopar } { x }

\cs_gset_protected:cx 2108 __cs_tmp:w { set_protected } { n }

\cs_gset_protected_nopar:cn 2109 __cs_tmp:w { set_protected } { x }

\cs_gset_protected_nopar:cx 2110 __cs_tmp:w { set_protected_nopar } { n }

\cs_gset_protected_nopar:cx 2111 __cs_tmp:w { set_protected_nopar } { x }

\cs_new:cn 2112 __cs_tmp:w { gset } { n }

\cs_new:cx 2113 __cs_tmp:w { gset } { x }

\cs_new_nopar:cn 2114 __cs_tmp:w { gset_nopar } { n }

\cs_new_nopar:cx 2115 __cs_tmp:w { gset_nopar } { x }

\cs_new_protected:cn 2116 __cs_tmp:w { gset_protected } { n }

\cs_new_protected:cx

\cs_new_protected_nopar:cn

\cs_new_protected_nopar:cx

```

2117 \__cs_tmp:w { gset_protected } { x }
2118 \__cs_tmp:w { gset_protected_nopar } { n }
2119 \__cs_tmp:w { gset_protected_nopar } { x }
2120 \__cs_tmp:w { new } { n }
2121 \__cs_tmp:w { new } { x }
2122 \__cs_tmp:w { new_nopar } { n }
2123 \__cs_tmp:w { new_nopar } { x }
2124 \__cs_tmp:w { new_protected } { n }
2125 \__cs_tmp:w { new_protected } { x }
2126 \__cs_tmp:w { new_protected_nopar } { n }
2127 \__cs_tmp:w { new_protected_nopar } { x }

```

(End definition for \cs_set:Nn. This function is documented on page 13.)

5.17 Checking control sequence equality

```

\cs_if_eq_p:NN Check if two control sequences are identical.
\cs_if_eq_p:cN 2128 \prg_new_conditional:Npnn \cs_if_eq:NN #1#2 { p , T , F , TF }
\cs_if_eq_p:Nc 2129 {
\cs_if_eq_p:cc 2130   \if_meaning:w #1#2
\cs_if_eq:NNTF 2131   \prg_return_true: \else: \prg_return_false: \fi:
\cs_if_eq:cNTF 2132 }
\cs_if_eq:NcTF 2133 \cs_new:Npn \cs_if_eq_p:cN { \exp_args:Nc \cs_if_eq_p:NN }
\cs_if_eq:NcTF 2134 \cs_new:Npn \cs_if_eq:cNTF { \exp_args:Nc \cs_if_eq:NNTF }
\cs_if_eq:NcTF 2135 \cs_new:Npn \cs_if_eq:cNT { \exp_args:Nc \cs_if_eq:NNT }
\cs_if_eq:NcTF 2136 \cs_new:Npn \cs_if_eq:cNF { \exp_args:Nc \cs_if_eq:NNF }
\cs_if_eq:NcTF 2137 \cs_new:Npn \cs_if_eq_p:Nc { \exp_args:NNc \cs_if_eq_p:NN }
\cs_if_eq:NcTF 2138 \cs_new:Npn \cs_if_eq:NcTF { \exp_args:NNc \cs_if_eq:NNTF }
\cs_if_eq:NcTF 2139 \cs_new:Npn \cs_if_eq:NcT { \exp_args:NNc \cs_if_eq:NNT }
\cs_if_eq:NcTF 2140 \cs_new:Npn \cs_if_eq:NcF { \exp_args:NNc \cs_if_eq:NNF }
\cs_if_eq:NcTF 2141 \cs_new:Npn \cs_if_eq_p:cc { \exp_args:Ncc \cs_if_eq_p:NN }
\cs_if_eq:NcTF 2142 \cs_new:Npn \cs_if_eq:ccTF { \exp_args:Ncc \cs_if_eq:NNTF }
\cs_if_eq:NcTF 2143 \cs_new:Npn \cs_if_eq:ccT { \exp_args:Ncc \cs_if_eq:NNT }
\cs_if_eq:NcTF 2144 \cs_new:Npn \cs_if_eq:ccF { \exp_args:Ncc \cs_if_eq:NNF }

```

(End definition for \cs_if_eq:NNTF. This function is documented on page 22.)

5.18 Diagnostic functions

```

2145 <@@=kernel>
\__kernel_chk_defined:NT Error if the variable #1 is not defined.
2146 \cs_new_protected:Npn \__kernel_chk_defined:NT #1#2
2147 {
2148   \cs_if_exist:NTF #1
2149   {#2}
2150   {
2151     \__kernel_msg_error:nxx { kernel } { variable-not-defined }
2152     { \token_to_str:N #1 }
2153   }
2154 }

```

(End definition for __kernel_chk_defined:NT.)

```

\__kernel_register_show:N
\__kernel_register_show:c
\__kernel_register_log:N
\__kernel_register_log:c
  \_kernel_register_show_aux:NN
  \_kernel_register_show_aux:nNN

```

Simply using the `\showthe` primitive does not allow for line-wrapping, so instead use `\tl_show:n` and `\tl_log:n` (defined in `l3tl` and that performs line-wrapping). This displays `>~⟨variable⟩=⟨value⟩`. We expand the value before-hand as otherwise some integers (such as `\currentgrouplevel` or `\currentgrouptype`) altered by the line-wrapping code would show wrong values.

```

2155 \cs_new_protected:Npn \__kernel_register_show:N
2156   { \__kernel_register_show_aux:NN \tl_show:n }
2157 \cs_new_protected:Npn \__kernel_register_show:c
2158   { \exp_args:Nc \__kernel_register_show:N }
2159 \cs_new_protected:Npn \__kernel_register_log:N
2160   { \__kernel_register_show_aux:NN \tl_log:n }
2161 \cs_new_protected:Npn \__kernel_register_log:c
2162   { \exp_args:Nc \__kernel_register_log:N }
2163 \cs_new_protected:Npn \__kernel_register_show_aux:NN #1#2
2164   {
2165     \__kernel_chk_defined:NT #2
2166     {
2167       \exp_args:No \__kernel_register_show_aux:nNN
2168       { \tex_the:D #2 } #2 #1
2169     }
2170   }
2171 \cs_new_protected:Npn \__kernel_register_show_aux:nNN #1#2#3
2172   { \exp_args:No #3 { \token_to_str:N #2 = #1 } }

```

(End definition for `__kernel_register_show:N` and others.)

```

\cs_show:N
\cs_show:c
\cs_log:N
\cs_log:c
\__kernel_show:NN

```

Some control sequences have a very long name or meaning. Thus, simply using TeX's primitive `\show` could lead to overlong lines. The output of this primitive is mimicked to some extent, then the re-built string is given to `\tl_show:n` or `\tl_log:n` for line-wrapping. We must expand the meaning before passing it to the wrapping code as otherwise we would wrongly see the definitions that are in place there. To get correct escape characters, set the `\escapechar` in a group; this also localizes the assignment performed by x-expansion. The `\cs_show:c` and `\cs_log:c` commands convert their argument to a control sequence within a group to avoid showing `\relax` for undefined control sequences.

```

2173 \cs_new_protected:Npn \cs_show:N { \__kernel_show:NN \tl_show:n }
2174 \cs_new_protected:Npn \cs_show:c
2175   { \group_begin: \exp_args:NNc \group_end: \cs_show:N }
2176 \cs_new_protected:Npn \cs_log:N { \__kernel_show:NN \tl_log:n }
2177 \cs_new_protected:Npn \cs_log:c
2178   { \group_begin: \exp_args:NNc \group_end: \cs_log:N }
2179 \cs_new_protected:Npn \__kernel_show:NN #1#2
2180   {
2181     \group_begin:
2182     \int_set:Nn \tex_escapechar:D { '\ }
2183     \exp_args:NNx
2184     \group_end:
2185     #1 { \token_to_str:N #2 = \cs_meaning:N #2 }
2186   }

```

(End definition for `\cs_show:N`, `\cs_log:N`, and `__kernel_show:NN`. These functions are documented on page 16.)

5.19 Decomposing a macro definition

`\cs_prefix_spec:N`
`\cs_argument_spec:N`
`\cs_replacement_spec:N`
`_kernel_prefix_arg_replacement:wN`

We sometimes want to test if a control sequence can be expanded to reveal a hidden value. However, we cannot just expand the macro blindly as it may have arguments and none might be present. Therefore we define these functions to pick either the prefix(es), the argument specification, or the replacement text from a macro. All of this information is returned as characters with catcode 12. If the token in question isn't a macro, the token `\scan_stop:` is returned instead.

```

2187 \use:x
2188 {
2189   \exp_not:n { \cs_new:Npn \_kernel_prefix_arg_replacement:wN #1 }
2190   \tl_to_str:n { macro : } \exp_not:n { #2 -> #3 \s__kernel_stop #4 }
2191 }
2192 { #4 {#1} {#2} {#3} }
2193 \cs_new:Npn \cs_prefix_spec:N #1
2194 {
2195   \token_if_macro:NTF #1
2196   {
2197     \exp_after:wN \_kernel_prefix_arg_replacement:wN
2198     \token_to_meaning:N #1 \s__kernel_stop \use_i:nnn
2199   }
2200   { \scan_stop: }
2201 }
2202 \cs_new:Npn \cs_argument_spec:N #1
2203 {
2204   \token_if_macro:NTF #1
2205   {
2206     \exp_after:wN \_kernel_prefix_arg_replacement:wN
2207     \token_to_meaning:N #1 \s__kernel_stop \use_ii:nnn
2208   }
2209   { \scan_stop: }
2210 }
2211 \cs_new:Npn \cs_replacement_spec:N #1
2212 {
2213   \token_if_macro:NTF #1
2214   {
2215     \exp_after:wN \_kernel_prefix_arg_replacement:wN
2216     \token_to_meaning:N #1 \s__kernel_stop \use_iii:nnn
2217   }
2218   { \scan_stop: }
2219 }
```

(End definition for `\cs_prefix_spec:N` and others. These functions are documented on page 18.)

5.20 Doing nothing functions

`\prg_do_nothing:`

This does not fit anywhere else!

```

2220 \cs_new:Npn \prg_do_nothing: { }
```

(End definition for `\prg_do_nothing:`. This function is documented on page 9.)

5.21 Breaking out of mapping functions

2221 `<@@=prg>`

`\prg_break_point:Nn` `\prg_map_break:Nn` In inline mappings, the nesting level must be reset at the end of the mapping, even when the user decides to break out. This is done by putting the code that must be performed as an argument of `__prg_break_point:Nn`. The breaking functions are then defined to jump to that point and perform the argument of `__prg_break_point:Nn`, before the user's code (if any). There is a check that we close the correct loop, otherwise we continue breaking.

```
2222 \cs_new_eq:NN \prg_break_point:Nn \use_ii:nn
2223 \cs_new:Npn \prg_map_break:Nn #1#2#3 \prg_break_point:Nn #4#5
2224 {
2225     #5
2226     \if_meaning:w #1 #4
2227     \exp_after:wN \use_iii:nnn
2228     \fi:
2229     \prg_map_break:Nn #1 {#2}
2230 }
```

(End definition for `\prg_break_point:Nn` and `\prg_map_break:Nn`. These functions are documented on page 113.)

`\prg_break_point:` `\prg_break:` `\prg_break:n` Very simple analogues of `\prg_break_point:Nn` and `\prg_map_break:Nn`, for use in fast short-term recursions which are not mappings, do not need to support nesting, and in which nothing has to be done at the end of the loop.

```
2231 \cs_new_eq:NN \prg_break_point: \prg_do_nothing:
2232 \cs_new:Npn \prg_break: #1 \prg_break_point: { }
2233 \cs_new:Npn \prg_break:n #1#2 \prg_break_point: {#1}
```

(End definition for `\prg_break_point:`, `\prg_break:`, and `\prg_break:n`. These functions are documented on page 114.)

5.22 Starting a paragraph

`\mode_leave_vertical:` The approach here is different to that used by L^AT_EX 2_ε or plain T_EX, which unbox a void box to force horizontal mode. That inserts the `\everypar` tokens *before* the re-inserted unboxing tokens. The approach here uses either the `\quitvmode` primitive or the equivalent protected macro. In vertical mode, the `\indent` primitive is inserted: this will switch to horizontal mode and insert `\everypar` tokens and nothing else. Unlike the L^AT_EX 2_ε version, the availability of ε-T_EX means using a mode test can be done at for example the start of an `\halign`.

```
2234 \cs_new_protected:Npn \mode_leave_vertical:
2235 {
2236     \if_mode_vertical:
2237     \exp_after:wN \tex_indent:D
2238     \fi:
2239 }
```

(End definition for `\mode_leave_vertical:`. This function is documented on page 24.)

2240 `</package>`

6 l3expan implementation

2241 `*package`

2242 `\@@=exp`

`\l__exp_internal_tl` The `\exp_` module has its private variable to temporarily store the result of `x`-type argument expansion. This is done to avoid interference with other functions using temporary variables.

(End definition for `\l__exp_internal_tl`.)

`\exp_after:wN` These are defined in `l3basics`, as they are needed “early”. This is just a reminder of that fact!

`\exp_not:N`

`\exp_not:n`

(End definition for `\exp_after:wN`, `\exp_not:N`, and `\exp_not:n`. These functions are documented on page 33.)

6.1 General expansion

In this section a general mechanism for defining functions that handle arguments is defined. These general expansion functions are expandable unless `x` is used. (Any version of `x` is going to have to use one of the L^AT_EX3 names for `\cs_set:Npx` at some point, and so is never going to be expandable.)

The definition of expansion functions with this technique happens in section 6.8. In section 6.2 some common cases are coded by a more direct method for efficiency, typically using calls to `\exp_after:wN`.

`\l__exp_internal_tl` This scratch token list variable is defined in `l3basics`.

(End definition for `\l__exp_internal_tl`.)

This code uses internal functions with names that start with `\::` to perform the expansions. All macros are `long` since the tokens undergoing expansion may be arbitrary user input.

An argument manipulator `\::⟨Z⟩` always has signature `#1\:::#2#3` where `#1` holds the remaining argument manipulations to be performed, `\:::` serves as an end marker for the list of manipulations, `#2` is the carried over result of the previous expansion steps and `#3` is the argument about to be processed. One exception to this rule is `\::p`, which has to grab an argument delimited by a left brace.

`__exp_arg_next:nnn` `#1` is the result of an expansion step, `#2` is the remaining argument manipulations and `#3` is the current result of the expansion chain. This auxiliary function moves `#1` back after `#3` in the input stream and checks if any expansion is left to be done by calling `#2`. In by far the most cases we need to add a set of braces to the result of an argument manipulation so it is more effective to do it directly here. Actually, so far only the `c` of the final argument manipulation variants does not require a set of braces.

2243 `\cs_new:Npn __exp_arg_next:nnn #1#2#3 { #2 \::: { #3 {#1} } }`

2244 `\cs_new:Npn __exp_arg_next:Nnn #1#2#3 { #2 \::: { #3 #1 } }`

(End definition for `__exp_arg_next:nnn` and `__exp_arg_next:Nnn`.)

`\:::` The end marker is just another name for the identity function.

2245 `\cs_new:Npn \::: #1 {#1}`

(End definition for `\:::`. This function is documented on page 37.)

- `\::n`** This function is used to skip an argument that doesn't need to be expanded.
- ```
2246 \cs_new:Npn \::n #1 \::: #2#3 { #1 \::: { #2 {#3} } }
```
- (End definition for `\::n`. This function is documented on page 37.)
- `\::N`** This function is used to skip an argument that consists of a single token and doesn't need to be expanded.
- ```
2247 \cs_new:Npn \::N #1 \::: #2#3 { #1 \::: {#2#3} }
```
- (End definition for `\::N`. This function is documented on page 37.)
- `\::p`** This function is used to skip an argument that is delimited by a left brace and doesn't need to be expanded. It is not wrapped in braces in the result.
- ```
2248 \cs_new:Npn \::p #1 \::: #2#3# { #1 \::: {#2#3} }
```
- (End definition for `\::p`. This function is documented on page 37.)
- `\::c`** This function is used to skip an argument that is turned into a control sequence without expansion.
- ```
2249 \cs_new:Npn \::c #1 \::: #2#3
2250 { \exp_after:wN \__exp_arg_next:Nnn \cs:w #3 \cs_end: {#1} {#2} }
```
- (End definition for `\::c`. This function is documented on page 37.)
- `\::o`** This function is used to expand an argument once.
- ```
2251 \cs_new:Npn \::o #1 \::: #2#3
2252 { \exp_after:wN __exp_arg_next:nnn \exp_after:wN {#3} {#1} {#2} }
```
- (End definition for `\::o`. This function is documented on page 37.)
- `\::e`** With the `\expanded` primitive available, just expand. Otherwise defer to `\exp_args:Ne` implemented later.
- ```
2253 \cs_if_exist:NTF \tex_expanded:D
2254 {
2255   \cs_new:Npn \::e #1 \::: #2#3
2256   { \tex_expanded:D { \exp_not:n { #1 \::: } { \exp_not:n {#2} {#3} } } }
2257 }
2258 {
2259   \cs_new:Npn \::e #1 \::: #2#3
2260   { \exp_args:Ne \__exp_arg_next:nnn {#3} {#1} {#2} }
2261 }
```
- (End definition for `\::e`. This function is documented on page 37.)
- `\::f`** This function is used to expand a token list until the first unexpandable token is found. This is achieved through `\exp:w \exp_end_continue_f:w` that expands everything in its way following it. This scanning procedure is terminated once the expansion hits something non-expandable (if that is a space it is removed). We introduce `\exp_stop_f:` to mark such an end-of-expansion marker. For example, f-expanding `\cs_set_eq:Nc \aaa { b \l_tmpa_tl b }` where `\l_tmpa_tl` contains the characters `lur` gives `\tex_let:D \aaa = \blurb` which then turns out to start with the non-expandable token `\tex_let:D`. Since the expansion of `\exp:w \exp_end_continue_f:w` is empty, we wind up with a fully expanded list, only `TEX` has not tried to execute any of

the non-expandable tokens. This is what differentiates this function from the `x` argument type.

```

2262 \cs_new:Npn \::f #1 \::: #2#3
2263 {
2264   \exp_after:wN \__exp_arg_next:nnn
2265   \exp_after:wN { \exp:w \exp_end_continue_f:w #3 }
2266   {#1} {#2}
2267 }
2268 \use:nn { \cs_new_eq:NN \exp_stop_f: } { ~ }
```

(End definition for `\::f` and `\exp_stop_f:`. These functions are documented on page 37.)

\::x This function is used to expand an argument fully. We build in the expansion of `__exp_arg_next:nnn`.

```

2269 \cs_new_protected:Npn \::x #1 \::: #2#3
2270 {
2271   \cs_set_nopar:Npx \l__exp_internal_tl
2272   { \exp_not:n { #1 \::: } { \exp_not:n {#2} {#3} } }
2273   \l__exp_internal_tl
2274 }
```

(End definition for `\::x`. This function is documented on page 37.)

\::v These functions return the value of a register, i.e., one of `tl`, `clist`, `int`, `skip`, `dim`, **\::V** `muskip`, or built-in `TeX` register. The `V` version expects a single token whereas `v` like `c` creates a `csname` from its argument given in braces and then evaluates it as if it was a `V`. The `\exp:w` sets off an expansion similar to an `f`-type expansion, which we terminate using `\exp_end:`. The argument is returned in braces.

```

2275 \cs_new:Npn \::V #1 \::: #2#3
2276 {
2277   \exp_after:wN \__exp_arg_next:nnn
2278   \exp_after:wN { \exp:w \__exp_eval_register:N #3 }
2279   {#1} {#2}
2280 }
2281 \cs_new:Npn \::v #1 \::: #2#3
2282 {
2283   \exp_after:wN \__exp_arg_next:nnn
2284   \exp_after:wN { \exp:w \__exp_eval_register:c {#3} }
2285   {#1} {#2}
2286 }
```

(End definition for `\::v` and `\::V`. These functions are documented on page 37.)

`__exp_eval_register:N` This function evaluates a register. Now a register might exist as one of two things: A parameter-less macro or a built-in `TeX` register such as `\count`. For the `TeX` registers we have to utilize a `\the` whereas for the macros we merely have to expand them once. The trick is to find out when to use `\the` and when not to. What we want here is to find out whether the token expands to something else when hit with `\exp_after:wN`. The technique is to compare the meaning of the token in question when it has been prefixed with `\exp_not:N` and the token itself. If it is a macro, the prefixed `\exp_not:N` temporarily turns it into the primitive `\scan_stop:`.

```

2287 \cs_new:Npn \__exp_eval_register:N #1
2288 {
2289   \exp_after:wN \if_meaning:w \exp_not:N #1 #1
```

If the token was not a macro it may be a malformed variable from a `c` expansion in which case it is equal to the primitive `\scan_stop:`. In that case we throw an error. We could let `TeX` do it for us but that would result in the rather obscure

```
! You can't use '\relax' after \the.
```

which while quite true doesn't give many hints as to what actually went wrong. We provide something more sensible.

```
2290     \if_meaning:w \scan_stop: #1
2291     \__exp_eval_error_msg:w
2292     \fi:
```

The next bit requires some explanation. The function must be initiated by `\exp:w` and we want to terminate this expansion chain by inserting the `\exp_end:` token. However, we have to expand the register `#1` before we do that. If it is a `TeX` register, we need to execute the sequence `\exp_after:wN \exp_end: \tex_the:D #1` and if it is a macro we need to execute `\exp_after:wN \exp_end: #1`. We therefore issue the longer of the two sequences and if the register is a macro, we remove the `\tex_the:D`.

```
2293     \else:
2294     \exp_after:wN \use_i_ii:nnn
2295     \fi:
2296     \exp_after:wN \exp_end: \tex_the:D #1
2297   }
2298 \cs_new:Npn \__exp_eval_register:c #1
2299 { \exp_after:wN \__exp_eval_register:N \cs:w #1 \cs_end: }
```

Clean up nicely, then call the undefined control sequence. The result is an error message looking like this:

```
! Undefined control sequence.
<argument> \LaTeX3 error:
                               Erroneous variable used!
1.55 \tl_set:Nv \l_tmpa_tl {undefined_tl}

2300 \cs_new:Npn \__exp_eval_error_msg:w #1 \tex_the:D #2
2301 {
2302   \fi:
2303   \fi:
2304   \__kernel_msg_expandable_error:nnn { kernel } { bad-variable } {#2}
2305   \exp_end:
2306 }
```

(End definition for `__exp_eval_register:N` and `__exp_eval_error_msg:w`.)

6.2 Hand-tuned definitions

One of the most important features of these functions is that they are fully expandable.

`\exp_args:Nc` In `l3basics`.

`\exp_args:cc` *(End definition for `\exp_args:Nc` and `\exp_args:cc`. These functions are documented on page 29.)*

`\exp_args:NNc` Here are the functions that turn their argument into csnames but are expandable.
`\exp_args:Ncc`
`\exp_args:Nccc`

```

2307 \cs_new:Npn \exp_args:NNc #1#2#3
2308 { \exp_after:wN #1 \exp_after:wN #2 \cs:w # 3\cs_end: }
2309 \cs_new:Npn \exp_args:Ncc #1#2#3
2310 { \exp_after:wN #1 \cs:w #2 \exp_after:wN \cs_end: \cs:w #3 \cs_end: }
2311 \cs_new:Npn \exp_args:Nccc #1#2#3#4
2312 {
2313   \exp_after:wN #1
2314   \cs:w #2 \exp_after:wN \cs_end:
2315   \cs:w #3 \exp_after:wN \cs_end:
2316   \cs:w #4 \cs_end:
2317 }
```

(End definition for `\exp_args:NNc`, `\exp_args:Ncc`, and `\exp_args:Nccc`. These functions are documented on page 31.)

`\exp_args:No` Those lovely runs of expansion!
`\exp_args:NNo`
`\exp_args:NNNo`

```

2318 \cs_new:Npn \exp_args:No #1#2 { \exp_after:wN #1 \exp_after:wN {#2} }
2319 \cs_new:Npn \exp_args:NNo #1#2#3
2320 { \exp_after:wN #1 \exp_after:wN #2 \exp_after:wN {#3} }
2321 \cs_new:Npn \exp_args:NNNo #1#2#3#4
2322 { \exp_after:wN #1 \exp_after:wN#2 \exp_after:wN #3 \exp_after:wN {#4} }
```

(End definition for `\exp_args:No`, `\exp_args:NNo`, and `\exp_args:NNNo`. These functions are documented on page 30.)

`\exp_args:Ne` When the `\expanded` primitive is available, use it. Otherwise use `__exp_e:nn`, defined later, to fully expand tokens.

```

2323 \cs_if_exist:NTF \tex_expanded:D
2324 {
2325   \cs_new:Npn \exp_args:Ne #1#2
2326   { \exp_after:wN #1 \tex_expanded:D { {#2} } }
2327 }
2328 {
2329   \cs_new:Npn \exp_args:Ne #1#2
2330   {
2331     \exp_after:wN #1 \exp_after:wN
2332     { \exp:w \__exp_e:nn {#2} { } }
2333   }
2334 }
```

(End definition for `\exp_args:Ne`. This function is documented on page 30.)

`\exp_args:Nf`
`\exp_args:Nv`
`\exp_args:Nv`

```

2335 \cs_new:Npn \exp_args:Nf #1#2
2336 { \exp_after:wN #1 \exp_after:wN { \exp:w \exp_end_continue_f:w #2 } }
2337 \cs_new:Npn \exp_args:Nv #1#2
2338 {
2339   \exp_after:wN #1 \exp_after:wN
2340   { \exp:w \__exp_eval_register:c {#2} }
2341 }
2342 \cs_new:Npn \exp_args:Nv #1#2
2343 {
2344   \exp_after:wN #1 \exp_after:wN
2345   { \exp:w \__exp_eval_register:N #2 }
2346 }
```

(End definition for `\exp_args:Nf`, `\exp_args:Nv`, and `\exp_args:Nv`. These functions are documented on page 30.)

`\exp_args:NNV` Some more hand-tuned function with three arguments. If we forced that an `o` argument always has braces, we could implement `\exp_args:Nco` with less tokens and only two arguments.

```

2347 \cs_new:Npn \exp_args:NNV #1#2#3
2348 {
2349   \exp_after:wN #1
2350   \exp_after:wN #2
2351   \exp_after:wN { \exp:w \__exp_eval_register:N #3 }
2352 }
2353 \cs_new:Npn \exp_args:NNv #1#2#3
2354 {
2355   \exp_after:wN #1
2356   \exp_after:wN #2
2357   \exp_after:wN { \exp:w \__exp_eval_register:c {#3} }
2358 }
2359 \cs_if_exist:NTF \tex_expanded:D
2360 {
2361   \cs_new:Npn \exp_args:NNe #1#2#3
2362   {
2363     \exp_after:wN #1
2364     \exp_after:wN #2
2365     \tex_expanded:D { {#3} }
2366   }
2367 }
2368 { \cs_new:Npn \exp_args:NNe { \::N \::e \::: } }
2369 \cs_new:Npn \exp_args:NNf #1#2#3
2370 {
2371   \exp_after:wN #1
2372   \exp_after:wN #2
2373   \exp_after:wN { \exp:w \exp_end_continue_f:w #3 }
2374 }
2375 \cs_new:Npn \exp_args:Nco #1#2#3
2376 {
2377   \exp_after:wN #1
2378   \cs:w #2 \exp_after:wN \cs_end:
2379   \exp_after:wN {#3}
2380 }
2381 \cs_new:Npn \exp_args:NcV #1#2#3
2382 {
2383   \exp_after:wN #1
2384   \cs:w #2 \exp_after:wN \cs_end:
2385   \exp_after:wN { \exp:w \__exp_eval_register:N #3 }
2386 }
2387 \cs_new:Npn \exp_args:Ncv #1#2#3
2388 {
2389   \exp_after:wN #1
2390   \cs:w #2 \exp_after:wN \cs_end:
2391   \exp_after:wN { \exp:w \__exp_eval_register:c {#3} }
2392 }
2393 \cs_new:Npn \exp_args:Ncf #1#2#3
2394 {

```

```

2395     \exp_after:wN #1
2396     \cs:w #2 \exp_after:wN \cs_end:
2397     \exp_after:wN { \exp:w \exp_end_continue_f:w #3 }
2398   }
2399 \cs_new:Npn \exp_args:NVV #1#2#3
2400 {
2401   \exp_after:wN #1
2402   \exp_after:wN { \exp:w \exp_after:wN
2403     \__exp_eval_register:N \exp_after:wN #2 \exp_after:wN }
2404   \exp_after:wN { \exp:w \__exp_eval_register:N #3 }
2405 }

```

(End definition for `\exp_args:NNV` and others. These functions are documented on page 31.)

`\exp_args:NNNV`
`\exp_args:NNNv`
`\exp_args:NcNc`
`\exp_args:NcNo`
`\exp_args:Ncco`

A few more that we can hand-tune.

```

2406 \cs_new:Npn \exp_args:NNNV #1#2#3#4
2407 {
2408   \exp_after:wN #1
2409   \exp_after:wN #2
2410   \exp_after:wN #3
2411   \exp_after:wN { \exp:w \__exp_eval_register:N #4 }
2412 }
2413 \cs_new:Npn \exp_args:NNNv #1#2#3#4
2414 {
2415   \exp_after:wN #1
2416   \exp_after:wN #2
2417   \exp_after:wN #3
2418   \exp_after:wN { \exp:w \__exp_eval_register:c {#4} }
2419 }
2420 \cs_new:Npn \exp_args:NcNc #1#2#3#4
2421 {
2422   \exp_after:wN #1
2423   \cs:w #2 \exp_after:wN \cs_end:
2424   \exp_after:wN #3
2425   \cs:w #4 \cs_end:
2426 }
2427 \cs_new:Npn \exp_args:NcNo #1#2#3#4
2428 {
2429   \exp_after:wN #1
2430   \cs:w #2 \exp_after:wN \cs_end:
2431   \exp_after:wN #3
2432   \exp_after:wN {#4}
2433 }
2434 \cs_new:Npn \exp_args:Ncco #1#2#3#4
2435 {
2436   \exp_after:wN #1
2437   \cs:w #2 \exp_after:wN \cs_end:
2438   \cs:w #3 \exp_after:wN \cs_end:
2439   \exp_after:wN {#4}
2440 }

```

(End definition for `\exp_args:NNNV` and others. These functions are documented on page 32.)

`\exp_args:Nx`

```

2441 \cs_new_protected:Npn \exp_args:Nx #1#2
2442 { \use:x { \exp_not:N #1 {#2} } }

```

(End definition for `\exp_args:Nx`. This function is documented on page 31.)

6.3 Last-unbraced versions

`__exp_arg_last_unbraced:nn` There are a few places where the last argument needs to be available unbraced. First some helper macros.

```

\::o_unbraced
\::V_unbraced
\::v_unbraced
\::e_unbraced
\::f_unbraced
\::x_unbraced
2443 \cs_new:Npn \__exp_arg_last_unbraced:nn #1#2 { #2#1 }
2444 \cs_new:Npn \::o_unbraced \::: #1#2
2445 { \exp_after:wN \__exp_arg_last_unbraced:nn \exp_after:wN {#2} {#1} }
2446 \cs_new:Npn \::V_unbraced \::: #1#2
2447 {
2448   \exp_after:wN \__exp_arg_last_unbraced:nn
2449   \exp_after:wN { \exp:w \__exp_eval_register:N #2 } {#1}
2450 }
2451 \cs_new:Npn \::v_unbraced \::: #1#2
2452 {
2453   \exp_after:wN \__exp_arg_last_unbraced:nn
2454   \exp_after:wN { \exp:w \__exp_eval_register:c {#2} } {#1}
2455 }
2456 \cs_if_exist:NTF \tex_expanded:D
2457 {
2458   \cs_new:Npn \::e_unbraced \::: #1#2
2459   { \tex_expanded:D { \exp_not:n {#1} #2 } }
2460 }
2461 {
2462   \cs_new:Npn \::e_unbraced \::: #1#2
2463   { \exp:w \__exp_e:nn {#2} {#1} }
2464 }
2465 \cs_new:Npn \::f_unbraced \::: #1#2
2466 {
2467   \exp_after:wN \__exp_arg_last_unbraced:nn
2468   \exp_after:wN { \exp:w \exp_end_continue_f:w #2 } {#1}
2469 }
2470 \cs_new_protected:Npn \::x_unbraced \::: #1#2
2471 {
2472   \cs_set_nopar:Npx \l__exp_internal_tl { \exp_not:n {#1} #2 }
2473   \l__exp_internal_tl
2474 }

```

(End definition for `__exp_arg_last_unbraced:nn` and others. These functions are documented on page 37.)

`\exp_last_unbraced:No` Now the business end: most of these are hand-tuned for speed, but the general system is in place.

```

\exp_last_unbraced:NV
\exp_last_unbraced:Nv
\exp_last_unbraced:Ne
\exp_last_unbraced:Nf
\exp_last_unbraced:NNo
\exp_last_unbraced:NNV
\exp_last_unbraced:NNf
\exp_last_unbraced:Nco
\exp_last_unbraced:NcV
\exp_last_unbraced:NNNo
\exp_last_unbraced:NNNV
\exp_last_unbraced:NNNf
\exp_last_unbraced:Nno
\exp_last_unbraced:Noo
\exp_last_unbraced:Nfo
\exp_last_unbraced:NnNo
\exp_last_unbraced:NNNNo
\exp_last_unbraced:NNNNf
2475 \cs_new:Npn \exp_last_unbraced:No #1#2 { \exp_after:wN #1 #2 }
2476 \cs_new:Npn \exp_last_unbraced:NV #1#2
2477 { \exp_after:wN #1 \exp:w \__exp_eval_register:N #2 }
2478 \cs_new:Npn \exp_last_unbraced:Nv #1#2
2479 { \exp_after:wN #1 \exp:w \__exp_eval_register:c {#2} }
2480 \cs_if_exist:NTF \tex_expanded:D
2481 {

```



```

2482 \cs_new:Npn \exp_last_unbraced:Ne #1#2
2483 { \exp_after:wN #1 \tex_expanded:D {#2} }
2484 }
2485 { \cs_new:Npn \exp_last_unbraced:Ne { \::e_unbraced \::: } }
2486 \cs_new:Npn \exp_last_unbraced:Nf #1#2
2487 { \exp_after:wN #1 \exp:w \exp_end_continue_f:w #2 }
2488 \cs_new:Npn \exp_last_unbraced:NNo #1#2#3
2489 { \exp_after:wN #1 \exp_after:wN #2 #3 }
2490 \cs_new:Npn \exp_last_unbraced:NNV #1#2#3
2491 {
2492 \exp_after:wN #1
2493 \exp_after:wN #2
2494 \exp:w \__exp_eval_register:N #3
2495 }
2496 \cs_new:Npn \exp_last_unbraced:NNf #1#2#3
2497 {
2498 \exp_after:wN #1
2499 \exp_after:wN #2
2500 \exp:w \exp_end_continue_f:w #3
2501 }
2502 \cs_new:Npn \exp_last_unbraced:Nco #1#2#3
2503 { \exp_after:wN #1 \cs:w #2 \exp_after:wN \cs_end: #3 }
2504 \cs_new:Npn \exp_last_unbraced:NcV #1#2#3
2505 {
2506 \exp_after:wN #1
2507 \cs:w #2 \exp_after:wN \cs_end:
2508 \exp:w \__exp_eval_register:N #3
2509 }
2510 \cs_new:Npn \exp_last_unbraced:NNNo #1#2#3#4
2511 { \exp_after:wN #1 \exp_after:wN #2 \exp_after:wN #3 #4 }
2512 \cs_new:Npn \exp_last_unbraced:NNNV #1#2#3#4
2513 {
2514 \exp_after:wN #1
2515 \exp_after:wN #2
2516 \exp_after:wN #3
2517 \exp:w \__exp_eval_register:N #4
2518 }
2519 \cs_new:Npn \exp_last_unbraced:NNNf #1#2#3#4
2520 {
2521 \exp_after:wN #1
2522 \exp_after:wN #2
2523 \exp_after:wN #3
2524 \exp:w \exp_end_continue_f:w #4
2525 }
2526 \cs_new:Npn \exp_last_unbraced:Nno { \::n \::o_unbraced \::: }
2527 \cs_new:Npn \exp_last_unbraced:Noo { \::o \::o_unbraced \::: }
2528 \cs_new:Npn \exp_last_unbraced:Nfo { \::f \::o_unbraced \::: }
2529 \cs_new:Npn \exp_last_unbraced:NnNo { \::n \::N \::o_unbraced \::: }
2530 \cs_new:Npn \exp_last_unbraced:NNNNo #1#2#3#4#5
2531 { \exp_after:wN #1 \exp_after:wN #2 \exp_after:wN #3 \exp_after:wN #4 #5 }
2532 \cs_new:Npn \exp_last_unbraced:NNNNf #1#2#3#4#5
2533 {
2534 \exp_after:wN #1
2535 \exp_after:wN #2

```

```

2536     \exp_after:wN #3
2537     \exp_after:wN #4
2538     \exp:w \exp_end_continue_f:w #5
2539   }
2540   \cs_new_protected:Npn \exp_last_unbraced:Nx { \::x_unbraced \::: }

```

(End definition for `\exp_last_unbraced:No` and others. These functions are documented on page 33.)

If #2 is a single token then this can be implemented as

```

\cs_new:Npn \exp_last_two_unbraced:Noo #1 #2 #3
{ \exp_after:wN \exp_after:wN \exp_after:wN #1 \exp_after:wN #2 #3 }

```

However, for robustness this is not suitable. Instead, a bit of a shuffle is used to ensure that #2 can be multiple tokens.

```

2541 \cs_new:Npn \exp_last_two_unbraced:Noo #1#2#3
2542   { \exp_after:wN \__exp_last_two_unbraced:noN \exp_after:wN {#3} {#2} #1 }
2543 \cs_new:Npn \__exp_last_two_unbraced:noN #1#2#3
2544   { \exp_after:wN #3 #2 #1 }

```

(End definition for `\exp_last_two_unbraced:Noo` and `__exp_last_two_unbraced:noN`. This function is documented on page 33.)

6.4 Preventing expansion

`__kernel_exp_not:w` At the kernel level, we need the primitive behaviour to allow expansion *before* the brace group.

```

2545 \cs_new_eq:NN \__kernel_exp_not:w \tex_unexpanded:D

```

(End definition for `__kernel_exp_not:w`.)

`\exp_not:c` All these except `\exp_not:c` call the kernel-internal `__kernel_exp_not:w` namely `\tex_unexpanded:D`.

```

\exp_not:o \tex_unexpanded:D.
\exp_not:e
\exp_not:f
\exp_not:V
\exp_not:v
2546 \cs_new:Npn \exp_not:c #1 { \exp_after:wN \exp_not:N \cs:w #1 \cs_end: }
2547 \cs_new:Npn \exp_not:o #1 { \__kernel_exp_not:w \exp_after:wN {#1} }
2548 \cs_if_exist:NTF \tex_expanded:D
2549   {
2550     \cs_new:Npn \exp_not:e #1
2551       { \__kernel_exp_not:w \tex_expanded:D { {#1} } }
2552   }
2553   {
2554     \cs_new:Npn \exp_not:e
2555       { \__kernel_exp_not:w \exp_args:Ne \prg_do_nothing: }
2556   }
2557 \cs_new:Npn \exp_not:f #1
2558   { \__kernel_exp_not:w \exp_after:wN { \exp:w \exp_end_continue_f:w #1 } }
2559 \cs_new:Npn \exp_not:V #1
2560   {
2561     \__kernel_exp_not:w \exp_after:wN
2562     { \exp:w \__exp_eval_register:N #1 }
2563   }
2564 \cs_new:Npn \exp_not:v #1
2565   {
2566     \__kernel_exp_not:w \exp_after:wN
2567     { \exp:w \__exp_eval_register:c {#1} }
2568   }

```

(End definition for `\exp_not:c` and others. These functions are documented on page 34.)

6.5 Controlled expansion

`\exp:w` To trigger a sequence of “arbitrarily” many expansions we need a method to invoke TeX’s expansion mechanism in such a way that (a) we are able to stop it in a controlled manner and (b) the result of what triggered the expansion in the first place is null, i.e., that we do not get any unwanted side effects. There aren’t that many possibilities in TeX; in fact the one explained below might well be the only one (as normally the result of expansion is not null).

The trick here is to make use of the fact that `\tex_romannumeral:D` expands the tokens following it when looking for a number and that its expansion is null if that number turns out to be zero or negative. So we use that to start the expansion sequence: `\exp:w` is set equal to `\tex_romannumeral:D` in `l3basics`. To stop the expansion sequence in a controlled way all we need to provide is a constant integer zero as part of expanded tokens. As this is an integer constant it immediately stops `\tex_romannumeral:D`’s search for a number. Again, the definition of `\exp_end:` as the integer constant zero is in `l3basics`. (Note that according to our specification all tokens we expand initiated by `\exp:w` are supposed to be expandable (as well as their replacement text in the expansion) so we will not encounter a “number” that actually result in a roman numeral being generated. Or if we do then the programmer made a mistake.)

If on the other hand we want to stop the initial expansion sequence but continue with an f-type expansion we provide the alphabetic constant `^^@` that also represents 0 but this time TeX’s syntax for a *number* continues searching for an optional space (and it continues expansion doing that) — see TeXbook page 269 for details.

```
2569 \group_begin:
2570 \tex_catcode:D ‘^^@ = 13
2571 \cs_new_protected:Npn \exp_end_continue_f:w { ‘^^@ }
```

If the above definition ever appears outside its proper context the active character `^^@` will be executed so we turn this into an error. The test for existence covers the (unlikely) case that some other code has already defined `^^@`: this is true for example for `xmLTex.tex`.

```
2572 \if_cs_exist:N ^^@
2573 \else:
2574 \cs_new:Npn ^^@
2575 { \__kernel_msg_expandable_error:nn { kernel } { bad-exp-end-f } }
2576 \fi:
```

The same but grabbing an argument to remove spaces and braces.

```
2577 \cs_new:Npn \exp_end_continue_f:nw #1 { ‘^^@ #1 }
2578 \group_end:
```

(End definition for `\exp:w` and others. These functions are documented on page 36.)

6.6 Emulating e-type expansion

When the `\expanded` primitive is available it is used to implement e-type expansion; otherwise we emulate it.

```
2579 \cs_if_exist:NF \tex_expanded:D
2580 {
```

`__exp_e:nn` Repeatedly expand tokens, keeping track of fully-expanded tokens in the second argument to `__exp_e:nn`; this function eventually calls `__exp_e_end:nn` to leave `\exp_end:` in the input stream, followed by the result of the expansion. There are many special cases: spaces, brace groups, `\noexpand`, `\unexpanded`, `\the`, `\primitive`. While we use brace tricks `\if_false: { \fi:`, the expansion of this function is always triggered by `\exp:w` so brace balance is eventually restored after that is hit with a single step of expansion. Otherwise we could not nest e-type expansions within each other.

```

2581 \cs_new:Npn \__exp_e:nn #1
2582 {
2583   \if_false: { \fi:
2584     \tl_if_head_is_N_type:nTF {#1}
2585     { \__exp_e:N }
2586     {
2587       \tl_if_head_is_group:nTF {#1}
2588       { \__exp_e_group:n }
2589       {
2590         \tl_if_empty:nTF {#1}
2591         { \exp_after:wN \__exp_e_end:nn }
2592         { \exp_after:wN \__exp_e_space:nn }
2593         \exp_after:wN { \if_false: } \fi:
2594       }
2595     }
2596   #1
2597 }
2598 }
2599 \cs_new:Npn \__exp_e_end:nn #1#2 { \exp_end: #2 }

```

(End definition for `__exp_e:nn` and `__exp_e_end:nn`.)

`__exp_e_space:nn` For an explicit space character, remove it by f-expansion and put it in the (future) output.

```

2600 \cs_new:Npn \__exp_e_space:nn #1#2
2601 { \exp_args:Nf \__exp_e:nn {#1} { #2 ~ } }

```

(End definition for `__exp_e_space:nn`.)

`__exp_e_group:n` For a group, expand its contents, wrap it in two pairs of braces, and call `__exp_e_put:nn`. This function places the first item (the double-brace wrapped result) into the output. Importantly, `\tl_head:n` works even if the input contains quarks.

```

2602 \cs_new:Npn \__exp_e_group:n #1
2603 {
2604   \exp_after:wN \__exp_e_put:nn
2605   \exp_after:wN { \exp_after:wN { \exp_after:wN {
2606     \exp:w \if_false: } \fi: \__exp_e:nn {#1} { } } }
2607 }
2608 \cs_new:Npn \__exp_e_put:nn #1
2609 {
2610   \exp_args:NNo \exp_args:No \__exp_e_put:nnn
2611   { \tl_head:n {#1} } {#1}
2612 }
2613 \cs_new:Npn \__exp_e_put:nnn #1#2#3
2614 { \exp_args:No \__exp_e:nn { \use_none:n #2 } { #3 #1 } }

```

(End definition for `__exp_e_group:n`, `__exp_e_put:nn`, and `__exp_e_put:nnn`.)

`__exp_e:N` For an N-type token, call `__exp_e:Nnn` with arguments the *⟨first token⟩*, the remain-
`__exp_e:Nnn` ing tokens to expand and what's already been expanded. If the *⟨first token⟩* is non-
`__exp_e_protected:Nnn` expandable, including `\protected` (`\long` or not) macros, it is put in the result by
`__exp_e_expandable:Nnn` `__exp_e_protected:Nnn`. The four special primitives `\unexpanded`, `\noexpand`, `\the`,
`\primitive` are detected; otherwise the token is expanded by `__exp_e_expandable:Nnn`.

```

2615 \cs_new:Npn \__exp_e:N #1
2616 {
2617   \exp_after:wN \__exp_e:Nnn
2618   \exp_after:wN #1
2619   \exp_after:wN { \if_false: } \fi:
2620 }
2621 \cs_new:Npn \__exp_e:Nnn #1
2622 {
2623   \if_case:w
2624     \exp_after:wN \if_meaning:w \exp_not:N #1 #1 1 ~ \fi:
2625     \token_if_protected_macro:NT #1 { 1 ~ }
2626     \token_if_protected_long_macro:NT #1 { 1 ~ }
2627     \if_meaning:w \exp_not:n #1 2 ~ \fi:
2628     \if_meaning:w \exp_not:N #1 3 ~ \fi:
2629     \if_meaning:w \tex_the:D #1 4 ~ \fi:
2630     \if_meaning:w \tex_primitive:D #1 5 ~ \fi:
2631     0 ~
2632     \exp_after:wN \__exp_e_expandable:Nnn
2633   \or: \exp_after:wN \__exp_e_protected:Nnn
2634   \or: \exp_after:wN \__exp_e_unexpanded:Nnn
2635   \or: \exp_after:wN \__exp_e_noexpand:Nnn
2636   \or: \exp_after:wN \__exp_e_the:Nnn
2637   \or: \exp_after:wN \__exp_e_primitive:Nnn
2638   \fi:
2639   #1
2640 }
2641 \cs_new:Npn \__exp_e_protected:Nnn #1#2#3
2642 { \__exp_e:nn {#2} { #3 #1 } }
2643 \cs_new:Npn \__exp_e_expandable:Nnn #1#2
2644 { \exp_args:No \__exp_e:nn { #1 #2 } }

```

(End definition for `__exp_e:N` and others.)

`__exp_e_primitive:Nnn` We don't try hard to make sensible error recovery since the error recovery of `\tex_`-
`__exp_e_primitive_aux:NNw` `primitive:D` when followed by something else than a primitive depends on the engine.
`__exp_e_primitive_aux:NNnn` The only valid case is when what follows is N-type. Then distinguish special primitives
`__exp_e_primitive_other:NNnn` `\unexpanded`, `\noexpand`, `\the`, `\primitive` from other primitives. In the "other" case,
`__exp_e_primitive_other_aux:nNNnn` the only reasonable way to check if the primitive that follows `\tex_primitive:D` is
expandable is to expand and compare the before-expansion and after-expansion results.
If they coincide then probably the primitive is non-expandable and should be put in the
output together with `\tex_primitive:D` (one can cook up contrived counter-examples
where the true `\expanded` would have an infinite loop), and otherwise one should continue
expanding.

```

2645 \cs_new:Npn \__exp_e_primitive:Nnn #1#2
2646 {
2647   \if_false: { \fi:
2648   \tl_if_head_is_N_type:nTF {#2}
2649     { \__exp_e_primitive_aux:NNw #1 }

```

```

2650         {
2651             \__kernel_msg_expandable_error:nnn { kernel } { e-type }
2652             { Missing~primitive~name }
2653             \__exp_e_primitive_aux:NNw #1 \c_empty_tl
2654         }
2655     #2
2656 }
2657 }
2658 \cs_new:Npn \__exp_e_primitive_aux:NNw #1#2
2659 {
2660     \exp_after:wN \__exp_e_primitive_aux:NNnn
2661     \exp_after:wN #1
2662     \exp_after:wN #2
2663     \exp_after:wN { \if_false: } \fi:
2664 }
2665 \cs_new:Npn \__exp_e_primitive_aux:NNnn #1#2
2666 {
2667     \exp_args:Nf \str_case_e:nnTF { \cs_to_str:N #2 }
2668     {
2669         { unexpanded } { \__exp_e_unexpanded:NNn \exp_not:n }
2670         { noexpand } { \__exp_e_noexpand:NNn \exp_not:N }
2671         { the } { \__exp_e_the:NNn \tex_the:D }
2672         {
2673             \sys_if_engine_xetex:T { pdf }
2674             \sys_if_engine luatex:T { pdf }
2675             primitive
2676         } { \__exp_e_primitive:NNn #1 }
2677     }
2678     { \__exp_e_primitive_other:NNnn #1 #2 }
2679 }
2680 \cs_new:Npn \__exp_e_primitive_other:NNnn #1#2#3
2681 {
2682     \exp_args:No \__exp_e_primitive_other_aux:nNNnn
2683     { #1 #2 #3 }
2684     #1 #2 {#3}
2685 }
2686 \cs_new:Npn \__exp_e_primitive_other_aux:nNNnn #1#2#3#4#5
2687 {
2688     \str_if_eq:nnTF {#1} { #2 #3 #4 }
2689     { \__exp_e:nn {#4} { #5 #2 #3 } }
2690     { \__exp_e:nn {#1} {#5} }
2691 }

```

(End definition for __exp_e_primitive:NNn and others.)

__exp_e_noexpand:NNn The \noexpand primitive has no effect when followed by a token that is not N-type; otherwise __exp_e_put:nn can grab the next token and put it in the result unchanged.

```

2692 \cs_new:Npn \__exp_e_noexpand:NNn #1#2
2693 {
2694     \tl_if_head_is_N_type:nTF {#2}
2695     { \__exp_e_put:nn } { \__exp_e:nn } {#2}
2696 }

```

(End definition for __exp_e_noexpand:NNn.)

`__exp_e_unexpanded:Nnn`
`__exp_e_unexpanded:nn`
`__exp_e_unexpanded:nN`
`__exp_e_unexpanded:N`

The `\unexpanded` primitive expands and ignores any space, `\scan_stop:`, or token affected by `\exp_not:N`, then expects a brace group. Since we only support brace-balanced token lists it is impossible to support the case where the argument of `\unexpanded` starts with an implicit brace. Even though we want to expand and ignore spaces we cannot blindly f-expand because tokens affected by `\exp_not:N` should be discarded without being expanded further.

As usual distinguish four cases: brace group (the normal case, where we just put the item in the result), space (just f-expand to remove the space), empty (an error), or N-type *token*. In the last case call `__exp_e_unexpanded:nN` triggered by an f-expansion. Having a non-expandable *token* after `\unexpanded` is an error (we recover by passing `{}` to `\unexpanded`; this is different from T_EX because the error recovery of `\unexpanded` changes the balance of braces), unless that *token* is `\scan_stop:` or a space (recall that we don't implement the case of an implicit begin-group token). An expandable *token* is instead expanded, unless it is `\noexpand`. The latter primitive can be followed by an expandable N-type token (removed), by a non-expandable one (kept and later causing an error), by a space (removed by f-expansion), or by a brace group or nothing (later causing an error).

```

2697 \cs_new:Npn \__exp_e_unexpanded:Nnn #1 { \__exp_e_unexpanded:nn }
2698 \cs_new:Npn \__exp_e_unexpanded:nn #1
2699 {
2700   \tl_if_head_is_N_type:nTF {#1}
2701   {
2702     \exp_args:Nf \__exp_e_unexpanded:nn
2703     { \__exp_e_unexpanded:nN {#1} #1 }
2704   }
2705   {
2706     \tl_if_head_is_group:nTF {#1}
2707     { \__exp_e_put:nn }
2708     {
2709       \tl_if_empty:nTF {#1}
2710       {
2711         \__kernel_msg_expandable_error:nnn
2712         { kernel } { e-type }
2713         { \unexpanded missing~brace }
2714         \__exp_e_end:nn
2715       }
2716       { \exp_args:Nf \__exp_e_unexpanded:nn }
2717     }
2718   } {#1}
2719 }
2720 }
2721 \cs_new:Npn \__exp_e_unexpanded:nN #1#2
2722 {
2723   \exp_after:wN \if_meaning:w \exp_not:N #2 #2
2724   \exp_after:wN \use_i:nn
2725   \else:
2726     \exp_after:wN \use_ii:nn
2727   \fi:
2728   {
2729     \token_if_eq_catcode:NNTF #2 \c_space_token
2730     { \exp_stop_f: }
2731     {

```

```

2732         \token_if_eq_meaning:NNTF #2 \scan_stop:
2733         { \exp_stop_f: }
2734         {
2735             \__kernel_msg_expandable_error:nnn
2736             { kernel } { e-type }
2737             { \unexpanded missing-brace }
2738             { }
2739         }
2740     }
2741 }
2742 {
2743     \token_if_eq_meaning:NNTF #2 \exp_not:N
2744     {
2745         \exp_args:No \tl_if_head_is_N_type:nT { \use_none:n #1 }
2746         { \__exp_e_unexpanded:N }
2747     }
2748     { \exp_after:wN \exp_stop_f: #2 }
2749 }
2750 }
2751 \cs_new:Npn \__exp_e_unexpanded:N #1
2752 {
2753     \exp_after:wN \if_meaning:w \exp_not:N #1 #1 \else:
2754     \exp_after:wN \use_i:nn
2755     \fi:
2756     \exp_stop_f: #1
2757 }

```

(End definition for `__exp_e_unexpanded:Nnn` and others.)

`__exp_e_the:Nnn`
`__exp_e_the:N`
`__exp_e_the_toks_reg:N`

Finally implement `\the`. Followed by anything other than an N-type *<token>* this causes an error (we just let T_EX make one), otherwise we test the *<token>*. If the *<token>* is expandable, expand it. Otherwise it could be any kind of register, or things like `\numexpr`, so there is no way to deal with all cases. Thankfully, only `\toks` data needs to be protected from expansion since everything else gives a string of characters. If the *<token>* is `\toks` we find a number and unpack using the `the_toks` functions. If it is a token register we unpack it in a brace group and call `__exp_e_put:nn` to move it to the result. Otherwise we unpack and continue expanding (useless but safe) since it is basically impossible to have a handle on where the result of `\the` ends.

```

2758     \cs_new:Npn \__exp_e_the:Nnn #1#2
2759     {
2760         \tl_if_head_is_N_type:nTF {#2}
2761         { \if_false: { \fi: \__exp_e_the:N #2 } }
2762         { \exp_args:No \__exp_e:nn { \tex_the:D #2 } }
2763     }
2764     \cs_new:Npn \__exp_e_the:N #1
2765     {
2766         \exp_after:wN \if_meaning:w \exp_not:N #1 #1
2767         \exp_after:wN \use_i:nn
2768         \else:
2769         \exp_after:wN \use_ii:nn
2770         \fi:
2771         {
2772             \if_meaning:w \tex_toks:D #1
2773             \exp_after:wN \__exp_e_the_toks:wnn \int_value:w

```



```

2774         \exp_after:wN \_\_exp_e_the_toks:n
2775         \exp_after:wN { \int_value:w \if_false: } \fi:
2776     \else:
2777         \_\_exp_e_if_toks_register:NTF #1
2778         { \exp_after:wN \_\_exp_e_the_toks_reg:N }
2779         {
2780             \exp_after:wN \_\_exp_e:nn \exp_after:wN {
2781                 \tex_the:D \if_false: } \fi:
2782         }
2783         \exp_after:wN #1
2784     \fi:
2785 }
2786 {
2787     \exp_after:wN \_\_exp_e_the:Nnn \exp_after:wN ?
2788     \exp_after:wN { \exp:w \if_false: } \fi:
2789     \exp_after:wN \exp_end: #1
2790 }
2791 }
2792 \cs_new:Npn \_\_exp_e_the_toks_reg:N #1
2793 {
2794     \exp_after:wN \_\_exp_e_put:nn \exp_after:wN {
2795         \exp_after:wN {
2796             \tex_the:D \if_false: } \fi: #1 }
2797 }

```

(End definition for `__exp_e_the:Nnn`, `__exp_e_the:N`, and `__exp_e_the_toks_reg:N`.)

`__exp_e_the_toks:wnn` The calling function has applied `\int_value:w` so we collect digits with `__exp_e_the_toks:n` (which gets the token list as an argument) and `__exp_e_the_toks:N` (which gets the first token in case it is N-type). The digits are themselves collected into an `\int_value:w` argument to `__exp_e_the_toks:wnn`. Then that function unpacks the `\toks<number>` into the result. We include `?` because `__exp_e_put:nnn` removes one item from its second argument. Note that our approach is rather crude: in cases like `\the\toks12~34` the first `\int_value:w` removes the space and we will incorrectly unpack the `\the\toks1234`.

```

2798 \cs_new:Npn \_\_exp_e_the_toks:wnn #1; #2
2799 {
2800     \exp_args:No \_\_exp_e_put:nnn
2801     { \tex_the:D \tex_toks:D #1 } { ? #2 }
2802 }
2803 \cs_new:Npn \_\_exp_e_the_toks:n #1
2804 {
2805     \tl_if_head_is_N_type:NTF {#1}
2806     { \exp_after:wN \_\_exp_e_the_toks:N \if_false: { \fi: #1 } }
2807     { ; {#1} }
2808 }
2809 \cs_new:Npn \_\_exp_e_the_toks:N #1
2810 {
2811     \if_int_compare:w 10 < 9 \token_to_str:N #1 \exp_stop_f:
2812     \exp_after:wN \use_i:nn
2813 }
2814 \else:
2815     \exp_after:wN \use_ii:nn
2816 \fi:
2817 {

```

```

2817         #1
2818         \exp_after:wN \_exp_e_the_toks:n
2819         \exp_after:wN { \if_false: } \fi:
2820     }
2821     {
2822         \exp_after:wN ;
2823         \exp_after:wN { \if_false: } \fi: #1
2824     }
2825 }

```

(End definition for `_exp_e_the_toks:wnn`, `_exp_e_the_toks:n`, and `_exp_e_the_toks:N`.)

`_exp_e_if_toks_register:N` We need to detect both `\toks` registers like `\toks@` in L^AT_EX 2_ε and parameters such as `\everypar`, as the result of unpacking the register should not expand further. Registers are found by `\token_if_toks_register:N` by inspecting the meaning. The list of parameters is finite so we just use a `\cs_if_exist:cTF` test to look up in a table. We abuse `\cs_to_str:N`'s ability to remove a leading escape character whatever it is.

```

\exp_e_the_everydisplay:
  \_exp_e_the_everyeof: 2826 \prg_new_conditional:Npnn \_exp_e_if_toks_register:N #1 { TF }
  \_exp_e_the_everyhbox: 2827 {
  \_exp_e_the_everyjob: 2828   \token_if_toks_register:N #1 { \prg_return_true: }
  \_exp_e_the_everymath: 2829   {
  \_exp_e_the_everypar: 2830     \cs_if_exist:cTF
  \_exp_e_the_everyvbox: 2831     {
  \_exp_e_the_output: 2832       \_exp_e_the_
  \_exp_e_the_pdfpageattr: 2833       \exp_after:wN \cs_to_str:N
  \_exp_e_the_pdfpageresources: 2834       \token_to_meaning:N #1
  \_exp_e_the_pdfpagesattr: 2835       :
  \_exp_e_the_pdfpkmode: 2836     } { \prg_return_true: } { \prg_return_false: }
  2837   }
  2838 }
2839 \cs_new_eq:NN \_exp_e_the_XeTeXinterchartoks: ?
2840 \cs_new_eq:NN \_exp_e_the_errhelp: ?
2841 \cs_new_eq:NN \_exp_e_the_everycr: ?
2842 \cs_new_eq:NN \_exp_e_the_everydisplay: ?
2843 \cs_new_eq:NN \_exp_e_the_everyeof: ?
2844 \cs_new_eq:NN \_exp_e_the_everyhbox: ?
2845 \cs_new_eq:NN \_exp_e_the_everyjob: ?
2846 \cs_new_eq:NN \_exp_e_the_everymath: ?
2847 \cs_new_eq:NN \_exp_e_the_everypar: ?
2848 \cs_new_eq:NN \_exp_e_the_everyvbox: ?
2849 \cs_new_eq:NN \_exp_e_the_output: ?
2850 \cs_new_eq:NN \_exp_e_the_pdfpageattr: ?
2851 \cs_new_eq:NN \_exp_e_the_pdfpageresources: ?
2852 \cs_new_eq:NN \_exp_e_the_pdfpagesattr: ?
2853 \cs_new_eq:NN \_exp_e_the_pdfpkmode: ?

```

(End definition for `_exp_e_if_toks_register:N` and others.)

We are done emulating e-type argument expansion when `\expanded` is unavailable.

```

2854 }

```

6.7 Defining function variants

```

2855 <@@=cs>

```

`\s__cs_mark` Internal scan marks. No l3quark yet, so do things by hand.

`\s__cs_stop` 2856 `\cs_new_eq:NN \s__cs_mark \scan_stop:`
 2857 `\cs_new_eq:NN \s__cs_stop \scan_stop:`

(End definition for `\s__cs_mark` and `\s__cs_stop`.)

`\q__cs_recursion_stop` Internal recursion quarks. No l3quark yet, so do things by hand.

2858 `\cs_new:Npn \q__cs_recursion_stop { \q__cs_recursion_stop }`

(End definition for `\q__cs_recursion_stop`.)

`__cs_use_none_delimit_by_s_stop:w` Internal scan marks.

`__cs_use_i_delimit_by_s_stop:nw` 2859 `\cs_new:Npn __cs_use_none_delimit_by_s_stop:w #1 \s__cs_stop { }`
`__cs_use_none_delimit_by_q_recursion_stop:w` 2860 `\cs_new:Npn __cs_use_i_delimit_by_s_stop:nw #1 #2 \s__cs_stop {#1}`
 2861 `\cs_new:Npn __cs_use_none_delimit_by_q_recursion_stop:w`
 2862 `#1 \q__cs_recursion_stop { }`

(End definition for `__cs_use_none_delimit_by_s_stop:w`, `__cs_use_i_delimit_by_s_stop:nw`, and `__cs_use_none_delimit_by_q_recursion_stop:w`.)

`\cs_generate_variant:Nn` #1 : Base form of a function; e.g., `\tl_set:Nn`

`\cs_generate_variant:cn` #2 : One or more variant argument specifiers; e.g., `{Nx,c,cx}`

After making sure that the base form exists, test whether it is protected or not and define `__cs_tmp:w` as either `\cs_new:Npx` or `\cs_new_protected:Npx`, which is then used to define all the variants (except those involving x-expansion, always protected). Split up the original base function only once, to grab its name and signature. Then we wish to iterate through the comma list of variant argument specifiers, which we first convert to a string: the reason is explained later.

2863 `\cs_new_protected:Npn \cs_generate_variant:Nn #1#2`
 2864 `{`
 2865 `__cs_generate_variant:N #1`
 2866 `\use:x`
 2867 `{`
 2868 `__cs_generate_variant:nnNN`
 2869 `\cs_split_function:N #1`
 2870 `\exp_not:N #1`
 2871 `\tl_to_str:n {#2} ,`
 2872 `\exp_not:N \scan_stop: ,`
 2873 `\exp_not:N \q__cs_recursion_stop`
 2874 `}`
 2875 `}`
 2876 `\cs_new_protected:Npn \cs_generate_variant:cn`
 2877 `{ \exp_args:Nc \cs_generate_variant:Nn }`

(End definition for `\cs_generate_variant:Nn`. This function is documented on page 27.)

`__cs_generate_variant:N`

The goal here is to pick up protected parent functions. There are four cases: the parent function can be a primitive or a macro, and can be expandable or not. For non-expandable primitives, all variants should be protected; skipping the `\else:` branch is safe because non-expandable primitives cannot be T_EX conditionals.

`__cs_generate_variant:ww`

`__cs_generate_variant:wwNw`

The other case where variants should be protected is when the parent function is a protected macro: then `protected` appears in the meaning before the first occurrence of `macro`. The `ww` auxiliary removes everything in the meaning string after the first `ma`. We use `ma` rather than the full `macro` because the meaning of the `\firstmark` primitive (and

four others) can contain an arbitrary string after a leading `firstmark:`. Then, look for `pr` in the part we extracted: no need to look for anything longer: the only strings we can have are an empty string, `\long`, `\protected`, `\protected\long`, `\first`, `\top`, `\bot`, `\splittop`, or `\splitbot`, with `\` replaced by the appropriate escape character. If `pr` appears in the part before `ma`, the first `\s__cs_mark` is taken as an argument of the `wwNw` auxiliary, and `#3` is `\cs_new_protected:Npx`, otherwise it is `\cs_new:Npx`.

```

2878 \cs_new_protected:Npx \__cs_generate_variant:N #1
2879 {
2880   \exp_not:N \exp_after:wN \exp_not:N \if_meaning:w
2881   \exp_not:N \exp_not:N #1 #1
2882   \cs_set_eq:NN \exp_not:N \__cs_tmp:w \cs_new_protected:Npx
2883   \exp_not:N \else:
2884   \exp_not:N \exp_after:wN \exp_not:N \__cs_generate_variant:ww
2885   \exp_not:N \token_to_meaning:N #1 \tl_to_str:n { ma }
2886   \s__cs_mark
2887   \s__cs_mark \cs_new_protected:Npx
2888   \tl_to_str:n { pr }
2889   \s__cs_mark \cs_new:Npx
2890   \s__cs_stop
2891   \exp_not:N \fi:
2892 }
2893 \exp_last_unbraced:NNNNo
2894 \cs_new_protected:Npn \__cs_generate_variant:ww
2895   #1 { \tl_to_str:n { ma } } #2 \s__cs_mark
2896   { \__cs_generate_variant:wwNw #1 }
2897 \exp_last_unbraced:NNNNo
2898 \cs_new_protected:Npn \__cs_generate_variant:wwNw
2899   #1 { \tl_to_str:n { pr } } #2 \s__cs_mark #3 #4 \s__cs_stop
2900   { \cs_set_eq:NN \__cs_tmp:w #3 }

```

(End definition for `__cs_generate_variant:N`, `__cs_generate_variant:ww`, and `__cs_generate_variant:wwNw`.)

`__cs_generate_variant:nnNN` #1 : Base name.
 #2 : Base signature.
 #3 : Boolean.
 #4 : Base function.

If the boolean is `\c_false_bool`, the base function has no colon and we abort with an error; otherwise, set off a loop through the desired variant forms. The original function is retained as `#4` for efficiency.

```

2901 \cs_new_protected:Npn \__cs_generate_variant:nnNN #1#2#3#4
2902 {
2903   \if_meaning:w \c_false_bool #3
2904   \__kernel_msg_error:nnx { kernel } { missing-colon }
2905   { \token_to_str:c {#1} }
2906   \exp_after:wN \__cs_use_none_delimit_by_q_recursion_stop:w
2907   \fi:
2908   \__cs_generate_variant:Nnnw #4 {#1}{#2}
2909 }

```

(End definition for `__cs_generate_variant:nnNN`.)

`__cs_generate_variant:Nnnw` #1 : Base function.
 #2 : Base name.

#3 : Base signature.

#4 : Beginning of variant signature.

First check whether to terminate the loop over variant forms. Then, for each variant form, construct a new function name using the original base name, the variant signature consisting of l letters and the last $k - l$ letters of the base signature (of length k). For example, for a base function `\prop_put:Nnn` which needs a `cV` variant form, we want the new signature to be `cVn`.

There are further subtleties:

- In `\cs_generate_variant:Nn \foo:nnTF {xxTF}`, we must define `\foo:xxTF` using `\exp_args:Nxx`, rather than a hypothetical `\exp_args:NxxTF`. Thus, we wish to trim a common trailing part from the base signature and the variant signature.
- In `\cs_generate_variant:Nn \foo:on {ox}`, the function `\foo:ox` must be defined using `\exp_args:Nnx`, not `\exp_args:Nox`, to avoid double `o` expansion.
- Lastly, `\cs_generate_variant:Nn \foo:on {xn}` must trigger an error, because we do not have a means to replace `o`-expansion by `x`-expansion. More generally, we can only convert `N` to `c`, or convert `n` to `V`, `v`, `o`, `f`, `x`.

All this boils down to a few rules. Only `n` and `N`-type arguments can be replaced by `\cs_generate_variant:Nn`. Other argument types are allowed to be passed unchanged from the base form to the variant: in the process they are changed to `n` except for `N` and `p`-type arguments. A common trailing part is ignored.

We compare the base and variant signatures one character at a time within `x`-expansion. The result is given to `__cs_generate_variant:wwNN` (defined later) in the form `<processed variant signature> \s__cs_mark <errors> \s__cs_stop <base function> <new function>`. If all went well, `<errors>` is empty; otherwise, it is a kernel error message and some clean-up code.

Note the space after #3 and after the following brace group. Those are ignored by `TeX` when fetching the last argument for `__cs_generate_variant_loop:nNwN`, but can be used as a delimiter for `__cs_generate_variant_loop_end:nwwwNNnn`.

```

2910 \cs_new_protected:Npn \__cs_generate_variant:Nnnw #1#2#3#4 ,
2911 {
2912   \if_meaning:w \scan_stop: #4
2913   \exp_after:wN \__cs_use_none_delimit_by_q_recursion_stop:w
2914   \fi:
2915   \use:x
2916   {
2917     \exp_not:N \__cs_generate_variant:wwNN
2918     \__cs_generate_variant_loop:nNwN { }
2919     #4
2920     \__cs_generate_variant_loop_end:nwwwNNnn
2921     \s__cs_mark
2922     #3 ~
2923     { ~ { } \fi: \__cs_generate_variant_loop_long:wNNnn } ~
2924     { }
2925     \s__cs_stop
2926     \exp_not:N #1 {#2} {#4}
2927   }
2928   \__cs_generate_variant:Nnnw #1 {#2} {#3}
2929 }
```

(End definition for `_cs_generate_variant:Nnnw`.)

<code>_cs_generate_variant_loop:nNwN</code>	#1 : Last few consecutive letters common between the base and variant (more precisely,
<code>_cs_generate_variant_loop_base:N</code>	<code>_cs_generate_variant_same:N</code> <i><letter></i> for each letter).
<code>_cs_generate_variant_loop_same:w</code>	#2 : Next variant letter.
<code>_cs_generate_variant_loop_end:nwwwNNnn</code>	#3 : Remainder of variant form.
<code>_cs_generate_variant_loop_long:wNNnn</code>	#4 : Next base letter.
<code>_cs_generate_variant_loop_invalid:NNwNNnn</code>	
<code>_cs_generate_variant_loop_special:NNwNNnn</code>	

The first argument is populated by `_cs_generate_variant_loop_same:w` when a variant letter and a base letter match. It is flushed into the input stream whenever the two letters are different: if the loop ends before, the argument is dropped, which means that trailing common letters are ignored.

The case where the two letters are different is only allowed if the base is N and the variant is o, V, v, f or x. Otherwise, call `_cs_generate_variant_loop_invalid:NNwNNnn` to remove the end of the loop, get arguments at the end of the loop, and place an appropriate error message as a second argument of `_cs_generate_variant:wwNN`. If the letters are distinct and the base letter is indeed n or N, leave in the input stream whatever argument #1 was collected, and the next variant letter #2, then loop by calling `_cs_generate_variant_loop:nNwN`.

The loop can stop in three ways.

- If the end of the variant form is encountered first, #2 is `_cs_generate_variant_loop_end:nwwwNNnn` (expanded by the conditional `\if:w`), which inserts some tokens to end the conditional; grabs the *<base name>* as #7, the *<variant signature>* #8, the *<next base letter>* #1 and the part #3 of the base signature that wasn't read yet; and combines those into the *<new function>* to be defined.
- If the end of the base form is encountered first, #4 is `~{} \fi:` which ends the conditional (with an empty expansion), followed by `_cs_generate_variant_loop_long:wNNnn`, which places an error as the second argument of `_cs_generate_variant:wwNN`.
- The loop can be interrupted early if the requested expansion is unavailable, namely when the variant and base letters differ and the base is not the right one (n or N to support the variant). In that case too an error is placed as the second argument of `_cs_generate_variant:wwNN`.

Note that if the variant form has the same length as the base form, #2 is as described in the first point, and #4 as described in the second point above. The `_cs_generate_variant_loop_end:nwwwNNnn` breaking function takes the empty brace group in #4 as its first argument: this empty brace group produces the correct signature for the full variant.

```

2930 \cs_new:Npn \_cs_generate_variant_loop:nNwN #1#2#3 \s_cs_mark #4
2931 {
2932   \if:w #2 #4
2933     \exp_after:wN \_cs_generate_variant_loop_same:w
2934   \else:
2935     \if:w #4 \_cs_generate_variant_loop_base:N #2 \else:
2936       \if:w 0
2937         \if:w N #4 \else: \if:w n #4 \else: 1 \fi: \fi:
2938         \if:w \scan_stop: \_cs_generate_variant_loop_base:N #2 1 \fi:
2939         0
2940       \_cs_generate_variant_loop_special:NNwNNnn #4#2

```

```

2941         \else:
2942             \__cs_generate_variant_loop_invalid:NNwNNnn #4#2
2943         \fi:
2944     \fi:
2945 \fi:
2946 #1
2947 \prg_do_nothing:
2948 #2
2949 \__cs_generate_variant_loop:nNwN { } #3 \s__cs_mark
2950 }
2951 \cs_new:Npn \__cs_generate_variant_loop_base:N #1
2952 {
2953     \if:w c #1 N \else:
2954         \if:w o #1 n \else:
2955             \if:w V #1 n \else:
2956                 \if:w v #1 n \else:
2957                     \if:w f #1 n \else:
2958                         \if:w e #1 n \else:
2959                             \if:w x #1 n \else:
2960                                 \if:w n #1 n \else:
2961                                     \if:w N #1 N \else:
2962                                         \scan_stop:
2963                                         \fi:
2964                                     \fi:
2965                                 \fi:
2966                             \fi:
2967                         \fi:
2968                     \fi:
2969                 \fi:
2970             \fi:
2971         \fi:
2972     }
2973 \cs_new:Npn \__cs_generate_variant_loop_same:w
2974     #1 \prg_do_nothing: #2#3#4
2975     { #3 { #1 \__cs_generate_variant_same:N #2 } }
2976 \cs_new:Npn \__cs_generate_variant_loop_end:nwwwNNnn
2977     #1#2 \s__cs_mark #3 ~ #4 \s__cs_stop #5#6#7#8
2978 {
2979     \scan_stop: \scan_stop: \fi:
2980     \s__cs_mark \s__cs_stop
2981     \exp_not:N #6
2982     \exp_not:c { #7 : #8 #1 #3 }
2983 }
2984 \cs_new:Npn \__cs_generate_variant_loop_long:wNNnn #1 \s__cs_stop #2#3#4#5
2985 {
2986     \exp_not:n
2987     {
2988         \s__cs_mark
2989         \__kernel_msg_error:nxxx { kernel } { variant-too-long }
2990         {#5} { \token_to_str:N #3 }
2991         \use_none:nnn
2992         \s__cs_stop
2993         #3
2994         #3

```

```

2995     }
2996   }
2997   \cs_new:Npn \__cs_generate_variant_loop_invalid:NNwNNnn
2998     #1#2 \fi: \fi: \fi: #3 \s__cs_stop #4#5#6#7
2999   {
3000     \fi: \fi: \fi:
3001     \exp_not:n
3002     {
3003       \s__cs_mark
3004       \__kernel_msg_error:nxxxxx { kernel } { invalid-variant }
3005       {#7} { \token_to_str:N #5 } {#1} {#2}
3006       \use_none:nnn
3007       \s__cs_stop
3008       #5
3009       #5
3010     }
3011   }
3012   \cs_new:Npn \__cs_generate_variant_loop_special:NNwNNnn
3013     #1#2#3 \s__cs_stop #4#5#6#7
3014   {
3015     #3 \s__cs_stop #4 #5 {#6} {#7}
3016     \exp_not:n
3017     {
3018       \__kernel_msg_error:nxxxxx
3019       { kernel } { deprecated-variant }
3020       {#7} { \token_to_str:N #5 } {#1} {#2}
3021     }
3022   }

```

(End definition for __cs_generate_variant_loop:nNwN and others.)

__cs_generate_variant_same:N

When the base and variant letters are identical, don't do any expansion. For most argument types, we can use the n-type no-expansion, but the N and p types require a slightly different behaviour with respect to braces. For V-type this function could output N to avoid adding useless braces but that is not a problem.

```

3023   \cs_new:Npn \__cs_generate_variant_same:N #1
3024   {
3025     \if:w N #1 #1 \else:
3026       \if:w p #1 #1 \else:
3027         \token_to_str:N n
3028         \if:w n #1 \else:
3029           \__cs_generate_variant_loop_special:NNwNNnn #1#1
3030         \fi:
3031       \fi:
3032     \fi:
3033   }

```

(End definition for __cs_generate_variant_same:N.)

__cs_generate_variant:wwNN

If the variant form has already been defined, log its existence (provided log-functions is active). Otherwise, make sure that the \exp_args:N #3 form is defined, and if it contains x, change __cs_tmp:w locally to \cs_new_protected:Npx. Then define the variant by combining the \exp_args:N #3 variant and the base function.

```

3034   \cs_new_protected:Npn \__cs_generate_variant:wwNN

```



```

3035     #1 \s__cs_mark #2 \s__cs_stop #3#4
3036   {
3037     #2
3038     \cs_if_free:NT #4
3039     {
3040       \group_begin:
3041         \__cs_generate_internal_variant:n {#1}
3042         \__cs_tmp:w #4 { \exp_not:c { exp_args:N #1 } \exp_not:N #3 }
3043       \group_end:
3044     }
3045   }

```

(End definition for __cs_generate_variant:wwNN.)

__cs_generate_internal_variant:n
__cs_generate_internal_variant_loop:n

First test for the presence of x (this is where working with strings makes our lives easier), as the result should be protected, and the next variant to be defined using that internal variant should be protected (done by setting __cs_tmp:w). Then call __cs_generate_internal_variant:NNn with arguments \cs_new_protected:cpn \use:x (for protected) or \cs_new:cpn \tex_expanded:D (expandable) and the signature. If p appears in the signature, or if the function to be defined is expandable and the primitive \expanded is not available, or if there are more than 8 arguments, call some fall-back code that just puts the appropriate \:: commands. Otherwise, call __cs_generate_internal_one_go:NNn to construct the \exp_args:N... function as a macro taking up to 9 arguments and expanding them using \use:x or \tex_expanded:D.

```

3046 \cs_new_protected:Npx \__cs_generate_internal_variant:n #1
3047   {
3048     \exp_not:N \__cs_generate_internal_variant:wwnNwn
3049     #1 \s__cs_mark
3050     { \cs_set_eq:NN \exp_not:N \__cs_tmp:w \cs_new_protected:Npx }
3051     \cs_new_protected:cpn
3052     \use:x
3053     \token_to_str:N x \s__cs_mark
3054     { }
3055     \cs_new:cpn
3056     \exp_not:N \tex_expanded:D
3057     \s__cs_stop
3058     {#1}
3059   }
3060 \exp_last_unbraced:NNNNo
3061 \cs_new_protected:Npn \__cs_generate_internal_variant:wwnNwn #1
3062   { \token_to_str:N x } #2 \s__cs_mark #3#4#5#6 \s__cs_stop #7
3063   {
3064     #3
3065     \cs_if_free:cT { exp_args:N #7 }
3066     { \__cs_generate_internal_variant:NNn #4 #5 {#7} }
3067   }
3068 \cs_set_protected:Npn \__cs_tmp:w #1
3069   {
3070     \cs_new_protected:Npn \__cs_generate_internal_variant:NNn ##1##2##3
3071     {
3072       \if_catcode:w X \use_none:nnnnnnnn ##3
3073       \prg_do_nothing: \prg_do_nothing: \prg_do_nothing:
3074       \prg_do_nothing: \prg_do_nothing: \prg_do_nothing:
3075       \prg_do_nothing: \prg_do_nothing: X

```

```

3076         \exp_after:wN \_cs_generate_internal_test:Nw \exp_after:wN ##2
3077     \else:
3078         \exp_after:wN \_cs_generate_internal_test_aux:w \exp_after:wN #1
3079     \fi:
3080     ##3
3081     \s__cs_mark
3082     {
3083         \use:x
3084         {
3085             ##1 { exp_args:N ##3 }
3086             { \_cs_generate_internal_variant_loop:n ##3 { : \use_i:nn } }
3087         }
3088     }
3089     #1
3090     \s__cs_mark
3091     { \exp_not:n { \_cs_generate_internal_one_go:NNn ##1 ##2 {##3} } }
3092     \s__cs_stop
3093 }
3094 \cs_new_protected:Npn \_cs_generate_internal_test_aux:w
3095     ##1 #1 ##2 \s__cs_mark ##3 ##4 \s__cs_stop {##3}
3096 \cs_if_exist:NTF \tex_expanded:D
3097 {
3098     \cs_new_eq:NN \_cs_generate_internal_test:Nw
3099     \_cs_generate_internal_test_aux:w
3100 }
3101 {
3102     \cs_new_protected:Npn \_cs_generate_internal_test:Nw ##1
3103     {
3104         \if_meaning:w \tex_expanded:D ##1
3105         \exp_after:wN \_cs_generate_internal_test_aux:w
3106         \exp_after:wN #1
3107         \else:
3108         \exp_after:wN \_cs_generate_internal_test_aux:w
3109         \fi:
3110     }
3111 }
3112 }
3113 \exp_args:No \_cs_tmp:w { \token_to_str:N p }
3114 \cs_new_protected:Npn \_cs_generate_internal_one_go:NNn #1#2#3
3115 {
3116     \_cs_generate_internal_loop:nwnnw
3117     { \exp_not:N ##1 } 1 . { } { }
3118     #3 { ? \_cs_generate_internal_end:w } X ;
3119     23456789 { ? \_cs_generate_internal_long:w } ;
3120     #1 #2 {##3}
3121 }
3122 \cs_new_protected:Npn \_cs_generate_internal_loop:nwnnw #1#2 . #3#4#5#6 ; #7
3123 {
3124     \use_none:n #5
3125     \use_none:n #7
3126     \cs_if_exist_use:cF { \_cs_generate_internal_#5:NN }
3127     { \_cs_generate_internal_other:NN }
3128     #5 #7
3129     #7 .

```

```

3130     { #3 #1 } { #4 ## #2 }
3131     #6 ;
3132 }
3133 \cs_new_protected:Npn \__cs_generate_internal_N:NN #1#2
3134 { \__cs_generate_internal_loop:nwnnw { \exp_not:N ###2 } }
3135 \cs_new_protected:Npn \__cs_generate_internal_c:NN #1#2
3136 { \exp_args:No \__cs_generate_internal_loop:nwnnw { \exp_not:c {###2} } }
3137 \cs_new_protected:Npn \__cs_generate_internal_n:NN #1#2
3138 { \__cs_generate_internal_loop:nwnnw { { \exp_not:n {###2} } } }
3139 \cs_new_protected:Npn \__cs_generate_internal_x:NN #1#2
3140 { \__cs_generate_internal_loop:nwnnw { {###2} } }
3141 \cs_new_protected:Npn \__cs_generate_internal_other:NN #1#2
3142 {
3143     \exp_args:No \__cs_generate_internal_loop:nwnnw
3144     {
3145         \exp_after:wN
3146         {
3147             \exp:w \exp_args:NNc \exp_after:wN \exp_end:
3148             { exp_not:#1 } {###2}
3149         }
3150     }
3151 }
3152 \cs_new_protected:Npn \__cs_generate_internal_end:w #1 . #2#3#4 ; #5 ; #6#7#8
3153 { #6 { exp_args:N #8 } #3 { #7 {#2} } }
3154 \cs_new_protected:Npn \__cs_generate_internal_long:w #1 N #2#3 . #4#5#6#
3155 {
3156     \exp_args:Nx \__cs_generate_internal_long:nnnNNn
3157     { \__cs_generate_internal_variant_loop:n #2 #6 { : \use_i:nn } }
3158     {#4} {#5}
3159 }
3160 \cs_new:Npn \__cs_generate_internal_long:nnnNNn #1#2#3#4 ; ; #5#6#7
3161 { #5 { exp_args:N #7 } #3 { #6 { \exp_not:n {#1} {#2} } } }

```

This command grabs char by char outputting \::#1 (not expanded further). We avoid tests by putting a trailing : \use_i:nn, which leaves \cs_end: and removes the looping macro. The colon is in fact also turned into \::: so that the required structure for \exp_args:N... commands is correctly terminated.

```

3162 \cs_new:Npn \__cs_generate_internal_variant_loop:n #1
3163 {
3164     \exp_after:wN \exp_not:N \cs:w :: #1 \cs_end:
3165     \__cs_generate_internal_variant_loop:n
3166 }

```

(End definition for __cs_generate_internal_variant:n and __cs_generate_internal_variant_loop:n.)

\prg_generate_conditional_variant:Nnn

```

\__cs_generate_variant:nnNnn 3167 \cs_new_protected:Npn \prg_generate_conditional_variant:Nnn #1
\__cs_generate_variant:w      3168 {
\__cs_generate_variant:n      3169     \use:x
    \__cs_generate_variant_p_form:nnn 3170     {
    \__cs_generate_variant_T_form:nnn 3171         \__cs_generate_variant:nnNnn
    \__cs_generate_variant_F_form:nnn 3172         \cs_split_function:N #1
    \__cs_generate_variant_TF_form:nnn 3173     }
3174 }

```

```

3175 \cs_new_protected:Npn \__cs_generate_variant:nnNnn #1#2#3#4#5
3176 {
3177   \if_meaning:w \c_false_bool #3
3178     \__kernel_msg_error:nnx { kernel } { missing-colon }
3179     { \token_to_str:c {#1} }
3180     \__cs_use_i_delimit_by_s_stop:nw
3181   \fi:
3182   \exp_after:wN \__cs_generate_variant:w
3183   \tl_to_str:n {#5} , \scan_stop: , \q__cs_recursion_stop
3184   \__cs_use_none_delimit_by_s_stop:w \s__cs_mark {#1} {#2} {#4} \s__cs_stop
3185 }
3186 \cs_new_protected:Npn \__cs_generate_variant:w
3187   #1 , #2 \s__cs_mark #3#4#5
3188 {
3189   \if_meaning:w \scan_stop: #1 \scan_stop:
3190     \if_meaning:w \q__cs_nil #1 \q__cs_nil
3191     \use_i:nnn
3192   \fi:
3193   \exp_after:wN \__cs_use_none_delimit_by_q_recursion_stop:w
3194   \else:
3195     \cs_if_exist_use:cTF { __cs_generate_variant_#1_form:nnn }
3196     { {#3} {#4} {#5} }
3197     {
3198       \__kernel_msg_error:nnxx
3199       { kernel } { conditional-form-unknown }
3200       {#1} { \token_to_str:c { #3 : #4 } }
3201     }
3202   \fi:
3203   \__cs_generate_variant:w #2 \s__cs_mark {#3} {#4} {#5}
3204 }
3205 \cs_new_protected:Npn \__cs_generate_variant_p_form:nnn #1#2
3206 { \cs_generate_variant:cn { #1_p : #2 } }
3207 \cs_new_protected:Npn \__cs_generate_variant_T_form:nnn #1#2
3208 { \cs_generate_variant:cn { #1 : #2 T } }
3209 \cs_new_protected:Npn \__cs_generate_variant_F_form:nnn #1#2
3210 { \cs_generate_variant:cn { #1 : #2 F } }
3211 \cs_new_protected:Npn \__cs_generate_variant_TF_form:nnn #1#2
3212 { \cs_generate_variant:cn { #1 : #2 TF } }

```

(End definition for \prg_generate_conditional_variant:Nnn and others. This function is documented on page 107.)

\exp_args_generate:n This function is not used in the kernel hence we can use functions that are defined in later modules. It also does not need to be fast so use inline mappings. For each requested variant we check that there are no characters besides NnpcofVvx, in particular that there are no spaces. Then we just call the internal function.

```

3213 \cs_new_protected:Npn \exp_args_generate:n #1
3214 {
3215   \exp_args:No \clist_map_inline:nn { \tl_to_str:n {#1} }
3216   {
3217     \str_map_inline:nn {##1}
3218     {
3219       \str_if_in:nnF { NnpcofVvx } {####1}
3220       {

```

```

3221         \__kernel_msg_error:nnnn { kernel } { invalid-exp-args }
3222         {####1} {##1}
3223         \str_map_break:n { \use_none:nn }
3224     }
3225 }
3226 \__cs_generate_internal_variant:n {##1}
3227 }
3228 }

```

(End definition for `\exp_args_generate:n`. This function is documented on page 269.)

6.8 Definitions with the automated technique

Some of these could be done more efficiently, but the complexity of coding then becomes an issue. Notice that the auto-generated functions actually take no arguments themselves.

Here are the actual function definitions, using the helper functions above. The group is used because `__cs_generate_internal_variant:n` redefines `__cs_tmp:w` locally.

```

\exp_args:Nnc \exp_args:Nno \exp_args:NnV \exp_args:Nnv \exp_args:Nne
\exp_args:Nnf \exp_args:Noc \exp_args:Noo \exp_args:Nof \exp_args:NVo
\exp_args:Nfo \exp_args:Nff \exp_args:Nee \exp_args:NNx \exp_args:Ncx
\exp_args:Nnx \exp_args:Nox \exp_args:Nxo \exp_args:Nxx
3229 \cs_set_protected:Npn \__cs_tmp:w #1
3230 {
3231     \group_begin:
3232         \exp_args:No \__cs_generate_internal_variant:n
3233         { \tl_to_str:n {#1} }
3234     \group_end:
3235 }
3236 \__cs_tmp:w { nc }
3237 \__cs_tmp:w { no }
3238 \__cs_tmp:w { nV }
3239 \__cs_tmp:w { nv }
3240 \__cs_tmp:w { ne }
3241 \__cs_tmp:w { nf }
3242 \__cs_tmp:w { oc }
3243 \__cs_tmp:w { oo }
3244 \__cs_tmp:w { of }
3245 \__cs_tmp:w { Vo }
3246 \__cs_tmp:w { fo }
3247 \__cs_tmp:w { ff }
3248 \__cs_tmp:w { ee }
3249 \__cs_tmp:w { Nx }
3250 \__cs_tmp:w { cx }
3251 \__cs_tmp:w { nx }
3252 \__cs_tmp:w { ox }
3253 \__cs_tmp:w { xo }
3254 \__cs_tmp:w { xx }

```

(End definition for `\exp_args:Nnc` and others. These functions are documented on page 31.)

```

\exp_args:NNcf \exp_args:NNno \exp_args:NNnV \exp_args:NNoo
\exp_args:NNVV \exp_args:Ncno \exp_args:NcnV \exp_args:Ncoo
\exp_args:NcVV \exp_args:Nnnc \exp_args:Nnno \exp_args:Nnnf
\exp_args:Nnff \exp_args:Nooo \exp_args:Noof \exp_args:Nffo
\exp_args:Neee

```

```

3261 \__cs_tmp:w { cnV }
3262 \__cs_tmp:w { coo }
3263 \__cs_tmp:w { cVV }
3264 \__cs_tmp:w { nnc }
3265 \__cs_tmp:w { nno }
3266 \__cs_tmp:w { nnf }
3267 \__cs_tmp:w { nff }
3268 \__cs_tmp:w { ooo }
3269 \__cs_tmp:w { oof }
3270 \__cs_tmp:w { ffo }
3271 \__cs_tmp:w { eee }
3272 \__cs_tmp:w { NNx }
3273 \__cs_tmp:w { Nnx }
3274 \__cs_tmp:w { Nox }
3275 \__cs_tmp:w { nnx }
3276 \__cs_tmp:w { nox }
3277 \__cs_tmp:w { ccx }
3278 \__cs_tmp:w { cnx }
3279 \__cs_tmp:w { oox }

```

(End definition for `\exp_args:NNcf` and others. These functions are documented on page 32.)

```

3280 \</package>

```

7 l3quark implementation

The following test files are used for this code: `m3quark001.lvt`.

```

3281 \*package>

```

7.1 Quarks

```

3282 \@@=quark>

```

`\quark_new:N` Allocate a new quark.

```

3283 \cs_new_protected:Npn \quark_new:N #1
3284 {
3285   \__kernel_chk_if_free_cs:N #1
3286   \cs_gset_nopar:Npn #1 {#1}
3287 }

```

(End definition for `\quark_new:N`. This function is documented on page 38.)

`\q_nil` Some “public” quarks. `\q_stop` is an “end of argument” marker, `\q_nil` is a empty value and `\q_no_value` marks an empty argument.

```

\q_no_value 3288 \quark_new:N \q_nil
\q_stop     3289 \quark_new:N \q_mark
            3290 \quark_new:N \q_no_value
            3291 \quark_new:N \q_stop

```

(End definition for `\q_nil` and others. These variables are documented on page 39.)

`\q_recursion_tail` Quarks for ending recursions. Only ever used there! `\q_recursion_tail` is appended to whatever list structure we are doing recursion on, meaning it is added as a proper list item with whatever list separator is in use. `\q_recursion_stop` is placed directly after the list.

```
3292 \quark_new:N \q_recursion_tail
3293 \quark_new:N \q_recursion_stop
```

(End definition for `\q_recursion_tail` and `\q_recursion_stop`. These variables are documented on page 39.)

`\s__quark` Private scan mark used in `l3quark`. We don't have `l3scan` yet, so we declare the scan mark here and add it to the scan mark pool later.

```
3294 \cs_new_eq:NN \s__quark \scan_stop:
```

(End definition for `\s__quark`.)

`\q__quark_nil` Private quark use for some tests.

```
3295 \quark_new:N \q__quark_nil
```

(End definition for `\q__quark_nil`.)

`\quark_if_recursion_tail_stop:N` When doing recursions, it is easy to spend a lot of time testing if the end marker has been found. To avoid this, a dedicated end marker is used each time a recursion is set up. Thus if the marker is found everything can be wrapper up and finished off. The simple case is when the test can guarantee that only a single token is being tested. In this case, there is just a dedicated copy of the standard quark test. Both a gobbling version and one inserting end code are provided.

```
3296 \cs_new:Npn \quark_if_recursion_tail_stop:N #1
3297 {
3298   \if_meaning:w \q_recursion_tail #1
3299   \exp_after:wN \use_none_delimit_by_q_recursion_stop:w
3300   \fi:
3301 }
3302 \cs_new:Npn \quark_if_recursion_tail_stop_do:Nn #1
3303 {
3304   \if_meaning:w \q_recursion_tail #1
3305   \exp_after:wN \use_i_delimit_by_q_recursion_stop:nw
3306   \else:
3307   \exp_after:wN \use_none:n
3308   \fi:
3309 }
```

(End definition for `\quark_if_recursion_tail_stop:N` and `\quark_if_recursion_tail_stop_do:Nn`. These functions are documented on page 40.)

`\quark_if_recursion_tail_stop:n` See `\quark_if_nil:nTF` for the details. Expanding `__quark_if_recursion_tail:w` once in front of the tokens chosen here gives an empty result if and only if `#1` is exactly `\q_recursion_tail`.

```
\quark_if_recursion_tail_stop:o
\quark_if_recursion_tail_stop_do:nn
\quark_if_recursion_tail_stop_do:on
\__quark_if_recursion_tail:w
3310 \cs_new:Npn \quark_if_recursion_tail_stop:n #1
3311 {
3312   \tl_if_empty:oTF
3313   { \__quark_if_recursion_tail:w {} #1 {} ?! \q_recursion_tail ??? }
3314   { \use_none_delimit_by_q_recursion_stop:w }
3315   { }
```

```

3316 }
3317 \cs_new:Npn \quark_if_recursion_tail_stop_do:nn #1
3318 {
3319   \tl_if_empty:oTF
3320   { \__quark_if_recursion_tail:w {} #1 {} ?! \q_recursion_tail ??? }
3321   { \use_i_delimit_by_q_recursion_stop:nw }
3322   { \use_none:n }
3323 }
3324 \cs_new:Npn \__quark_if_recursion_tail:w
3325   #1 \q_recursion_tail #2 ? #3 ?! { #1 #2 }
3326 \cs_generate_variant:Nn \quark_if_recursion_tail_stop:n { o }
3327 \cs_generate_variant:Nn \quark_if_recursion_tail_stop_do:nn { o }

```

(End definition for `\quark_if_recursion_tail_stop:n`, `\quark_if_recursion_tail_stop_do:nn`, and `__quark_if_recursion_tail:w`. These functions are documented on page 40.)

`\quark_if_recursion_tail_break:NN`
`\quark_if_recursion_tail_break:nN`

Analogues of the `\quark_if_recursion_tail_stop...` functions. Break the mapping using #2.

```

3328 \cs_new:Npn \quark_if_recursion_tail_break:NN #1#2
3329 {
3330   \if_meaning:w \q_recursion_tail #1
3331   \exp_after:wN #2
3332   \fi:
3333 }
3334 \cs_new:Npn \quark_if_recursion_tail_break:nN #1#2
3335 {
3336   \tl_if_empty:oT
3337   { \__quark_if_recursion_tail:w {} #1 {} ?! \q_recursion_tail ??? }
3338   {#2}
3339 }

```

(End definition for `\quark_if_recursion_tail_break:NN` and `\quark_if_recursion_tail_break:nN`. These functions are documented on page 40.)

`\quark_if_nil_p:N`
`\quark_if_nil:N \underline{TF}`
`\quark_if_no_value_p:N`
`\quark_if_no_value_p:c`
`\quark_if_no_value:N \underline{TF}`
`\quark_if_no_value:c \underline{TF}`

Here we test if we found a special quark as the first argument. We better start with `\q_no_value` as the first argument since the whole thing may otherwise loop if #1 is wrongly given a string like `aabc` instead of a single token.⁸

```

3340 \prg_new_conditional:Npnn \quark_if_nil:N #1 { p, T, F, TF }
3341 {
3342   \if_meaning:w \q_nil #1
3343   \prg_return_true:
3344   \else:
3345     \prg_return_false:
3346   \fi:
3347 }
3348 \prg_new_conditional:Npnn \quark_if_no_value:N #1 { p, T, F, TF }
3349 {
3350   \if_meaning:w \q_no_value #1
3351   \prg_return_true:
3352   \else:
3353     \prg_return_false:
3354   \fi:
3355 }

```

⁸It may still loop in special circumstances however!


```

3356 \prg_generate_conditional_variant:Nnn \quark_if_no_value:N
3357 { c } { p , T , F , TF }

```

(End definition for \quark_if_nil:N_{TF} and \quark_if_no_value:N_{TF}. These functions are documented on page 39.)

\quark_if_nil_p:n Let us explain \quark_if_nil:n(TF). Expanding __quark_if_nil:w once is safe thanks to the trailing \q_nil ??!. The result of expanding once is empty if and only if both delimited arguments #1 and #2 are empty and #3 is delimited by the last tokens ?!. Thanks to the leading {}, the argument #1 is empty if and only if the argument of \quark_if_nil:n starts with \q_nil. The argument #2 is empty if and only if this \q_nil is followed immediately by ? or by {}?, coming either from the trailing tokens in the definition of \quark_if_nil:n, or from its argument. In the first case, __quark_if_nil:w is followed by {} \q_nil {}? ! \q_nil ??!, hence #3 is delimited by the final ?!, and the test returns true as wanted. In the second case, the result is not empty since the first ?! in the definition of \quark_if_nil:n stop #3. The auxiliary here is the same as __tl_if_empty_if:o, with the same comments applying.

```

3358 \prg_new_conditional:Npnn \quark_if_nil:n #1 { p , T , F , TF }
3359 {
3360   \__quark_if_empty_if:o
3361   { \__quark_if_nil:w {} #1 {} ? ! \q_nil ? ? ! }
3362   \prg_return_true:
3363   \else:
3364     \prg_return_false:
3365   \fi:
3366 }
3367 \cs_new:Npn \__quark_if_nil:w #1 \q_nil #2 ? #3 ? ! { #1 #2 }
3368 \prg_new_conditional:Npnn \quark_if_no_value:n #1 { p , T , F , TF }
3369 {
3370   \__quark_if_empty_if:o
3371   { \__quark_if_no_value:w {} #1 {} ? ! \q_no_value ? ? ! }
3372   \prg_return_true:
3373   \else:
3374     \prg_return_false:
3375   \fi:
3376 }
3377 \cs_new:Npn \__quark_if_no_value:w #1 \q_no_value #2 ? #3 ? ! { #1 #2 }
3378 \prg_generate_conditional_variant:Nnn \quark_if_nil:n
3379 { V , o } { p , TF , T , F }
3380 \cs_new:Npn \__quark_if_empty_if:o #1
3381 {
3382   \exp_after:wN \if_meaning:w \exp_after:wN \q_nil
3383   \__kernel_tl_to_str:w \exp_after:wN {#1} \q_nil
3384 }

```

(End definition for \quark_if_nil:n_{TF} and others. These functions are documented on page 39.)

__kernel_quark_new_test:N The function __kernel_quark_new_test:N defines #1 in a similar way as \quark_if_recursion_tail... functions (as described below), using \q_<namespace>-recursion_tail as the test quark and \q_<namespace>-recursion_stop as the delimiter quark, where the <namespace> is determined as the first _-delimited part in #1.

There are six possible function types which this function can define, and which is defined depends on the signature of the function being defined:

:n gives an analogue of \quark_if_recursion_tail_stop:n
:nn gives an analogue of \quark_if_recursion_tail_stop_do:nn
:nN gives an analogue of \quark_if_recursion_tail_break:nN
:N gives an analogue of \quark_if_recursion_tail_stop:N
:Nn gives an analogue of \quark_if_recursion_tail_stop_do:Nn
:NN gives an analogue of \quark_if_recursion_tail_break:NN

Any other signature causes an error, as does a function without signature.

Similar to __kernel_quark_new_test:N, but defines quark branching conditionals like __kernel_quark_new_conditional:Nn \quark_if_nil:nTF that test for the quark \q_<namespace>_<name>. The <namespace> and <name> are determined from the conditional #1, which must take the rather rigid form __<namespace>_quark_if_<name>:<arg spec>. There are only two cases for the <arg spec> here:

:n gives an analogue of \quark_if_nil:n(TF)
:N gives an analogue of \quark_if_nil:N(TF)

Any other signature causes an error, as does a function without signature. We use low-level emptiness tests as l3tl is not available yet when these functions are used; thankfully we only care about whether strings are empty so a simple \if_meaning:w \q_nil <string> \q_nil suffices.

```

3385 \cs_new_protected:Npn \__kernel_quark_new_test:N #1
3386 { \__quark_new_test_aux:Nx #1 { \__quark_module_name:N #1 } }
3387 \cs_new_protected:Npn \__quark_new_test_aux:Nn #1 #2
3388 {
3389   \if_meaning:w \q_nil #2 \q_nil
3390     \__kernel_msg_error:nxx { kernel } { invalid-quark-function }
3391     { \token_to_str:N #1 }
3392   \else:
3393     \__quark_new_test:Nccn #1
3394     { q_#2_recursion_tail } { q_#2_recursion_stop } { __#2 }
3395   \fi:
3396 }
3397 \cs_generate_variant:Nn \__quark_new_test_aux:Nn { Nx }
3398 \cs_new_protected:Npn \__quark_new_test:NNNn #1
3399 {
3400   \exp_last_unbraced:Nf \__quark_new_test_aux:nnNNnnnn
3401   { \cs_split_function:N #1 }
3402   #1 { test }
3403 }
3404 \cs_generate_variant:Nn \__quark_new_test:NNNn { Ncc }
3405 \cs_new_protected:Npn \__kernel_quark_new_conditional:Nn #1
3406 {
3407   \__quark_new_conditional:Nxxn #1
3408   { \__quark_quark_conditional_name:N #1 }
3409   { \__quark_module_name:N #1 }

```

```

3410 }
3411 \cs_new_protected:Npn \__quark_new_conditional:Nnnn #1#2#3#4
3412 {
3413   \if_meaning:w \q_nil #2 \q_nil
3414     \__kernel_msg_error:nxx { kernel } { invalid-quark-function }
3415     { \token_to_str:N #1 }
3416   \else:
3417     \if_meaning:w \q_nil #3 \q_nil
3418       \__kernel_msg_error:nxx { kernel } { invalid-quark-function }
3419       { \token_to_str:N #1 }
3420     \else:
3421       \exp_last_unbraced:Nf \__quark_new_test_aux:nnNNnnnn
3422       { \cs_split_function:N #1 }
3423       #1 { conditional }
3424       {#2} {#3} {#4}
3425     \fi:
3426   \fi:
3427 }
3428 \cs_generate_variant:Nn \__quark_new_conditional:Nnnn { Nxx }
3429 \cs_new_protected:Npn \__quark_new_test_aux:nnNNnnnn #1 #2 #3 #4 #5
3430 {
3431   \cs_if_exist_use:cTF { __quark_new_#5_#2:Nnnn } { #4 }
3432   {
3433     \__kernel_msg_error:nxxx { kernel } { invalid-quark-function }
3434     { \token_to_str:N #4 } {#2}
3435     \use_none:nnn
3436   }
3437 }

```

(End definition for __kernel_quark_new_test:N and others.)

```

\__quark_new_test_n:Nnnn
\__quark_new_test_nn:Nnnn
\__quark_new_test_N:Nnnn
\__quark_new_test_Nn:Nnnn
\__quark_new_test_NN:Nnnn
\__quark_new_test_NN:Nnnn

```

These macros implement the six possibilities mentioned above, passing the right arguments to __quark_new_test_aux_do:nnNNnnnnNNn, which defines some auxiliaries, and then to __quark_new_test_define_tl:nNnNNn (:n(n) variants) or to __quark_new_test_define_ifx:nNnNNn (:N(n)) which define the main conditionals.

```

3438 \cs_new_protected:Npn \__quark_new_test_n:Nnnn #1 #2 #3 #4
3439 {
3440   \__quark_new_test_aux_do:nnNNnnnnNNn {#4} #2 #3 { none } { } { } { }
3441   \__quark_new_test_define_tl:nNnNNn #1 { }
3442 }
3443 \cs_new_protected:Npn \__quark_new_test_nn:Nnnn #1 #2 #3 #4
3444 {
3445   \__quark_new_test_aux_do:nnNNnnnnNNn {#4} #2 #3 { i } { n } {##1} {##2}
3446   \__quark_new_test_define_tl:nNnNNn #1 { \use_none:n }
3447 }
3448 \cs_new_protected:Npn \__quark_new_test_nN:Nnnn #1 #2 #3 #4
3449 {
3450   \__quark_new_test_aux_do:nnNNnnnnNNn {#4} #2 #3 { i } { n } {##1} {##2}
3451   \__quark_new_test_define_break_tl:nNNNNn #1 { }
3452 }
3453 \cs_new_protected:Npn \__quark_new_test_N:Nnnn #1 #2 #3 #4
3454 {
3455   \__quark_new_test_aux_do:nnNNnnnnNNn {#4} #2 #3 { none } { } { } { }
3456   \__quark_new_test_define_ifx:nNnNNn #1 { }

```

```

3457 }
3458 \cs_new_protected:Npn \__quark_new_test_Nn:Nnnn #1 #2 #3 #4
3459 {
3460   \__quark_new_test_aux_do:nNNnnnnNNn {#4} #2 #3 { i } { n } {##1} {##2}
3461   \__quark_new_test_define_ifx:nNnNNn #1
3462   { \else: \exp_after:wN \use_none:n }
3463 }
3464 \cs_new_protected:Npn \__quark_new_test_NN:Nnnn #1 #2 #3 #4
3465 {
3466   \__quark_new_test_aux_do:nNNnnnnNNn {#4} #2 #3 { i } { n } {##1} {##2}
3467   \__quark_new_test_define_break_ifx:nNNNNn #1 { }
3468 }

```

(End definition for __quark_new_test_n:Nnnn and others.)

__quark_new_test_aux_do:nNNnnnnNNn __quark_new_test_aux_do:nNNnnnnNNn makes the control sequence names which will be used by __quark_test_define_aux:NNNNnnNNn, and then later by __quark_new_test_define_tl:nNnNNn or __quark_new_test_define_ifx:nNnNNn. The control sequences defined here are analogous to __quark_if_recursion_tail:w and to \use_-(none|i)_delimit_by_q_recursion_stop:(|n)w.

The name is composed by the name-space and the name of the quarks. Suppose __kernel_quark_new_test:N was used with:

```
\__kernel_quark_new_test:N \__test_quark_tail:n
```

then the first auxiliary will be __test_quark_recursion_tail:w, and the second one will be __test_use_none_delimit_by_q_recursion_stop:w.

Note that the actual quarks are *not* defined here. They should be defined separately using \quark_new:N.

```

3469 \cs_new_protected:Npn \__quark_new_test_aux_do:nNNnnnnNNn #1 #2 #3 #4 #5
3470 {
3471   \exp_args:Ncc \__quark_test_define_aux:NNNNnnNNn
3472   { #1 _quark_recursion_tail:w }
3473   { #1 _use_ #4 _delimit_by_q_recursion_stop: #5 w }
3474   #2 #3
3475 }
3476 \cs_new_protected:Npn \__quark_test_define_aux:NNNNnnNNn #1 #2 #3 #4 #5 #6 #7
3477 {
3478   \cs_gset:Npn #1 ##1 #3 ##2 ? ##3 ?! { ##1 ##2 }
3479   \cs_gset:Npn #2 ##1 #6 #4 {#5}
3480   #7 {##1} #1 #2 #3
3481 }

```

(End definition for __quark_new_test_aux_do:nNNnnnnNNn and __quark_test_define_aux:NNNNnnNNn.)

```

\__quark_new_test_define_tl:nNnNNn
\__quark_new_test_define_ifx:nNnNNn
\__quark_new_test_define_break_tl:nNNNNn
\__quark_new_test_define_break_ifx:nNNNNn

```

Finally, these two macros define the main conditional function using what's been set up before.

```

3482 \cs_new_protected:Npn \__quark_new_test_define_tl:nNnNNn #1 #2 #3 #4 #5 #6
3483 {
3484   \cs_new:Npn #5 #1
3485   {
3486     \tl_if_empty:oTF
3487     { #2 {} ##1 {} ?! #4 ??! }
3488     {#3} {#6}

```

```

3489     }
3490   }
3491   \cs_new_protected:Npn \__quark_new_test_define_ifx:nNnNnNn #1 #2 #3 #4 #5 #6
3492   {
3493     \cs_new:Npn #5 #1
3494     {
3495       \if_meaning:w #4 ##1
3496       \exp_after:wN #3
3497       #6
3498     \fi:
3499   }
3500 }
3501 \cs_new_protected:Npn \__quark_new_test_define_break_tl:nNNNNn #1 #2 #3
3502 { \__quark_new_test_define_tl:nNnNnNn {##1##2} #2 {##2} }
3503 \cs_new_protected:Npn \__quark_new_test_define_break_ifx:nNNNNNn #1 #2 #3
3504 { \__quark_new_test_define_ifx:nNnNnNn {##1##2} #2 {##2} }

```

(End definition for __quark_new_test_define_tl:nNnNnNn and others.)

__quark_new_conditional_n:Nnnn These macros implement the two possibilities for branching quark conditionals, passing the right arguments to __quark_new_conditional_aux_do:NNnnnn, which defines some auxiliaries and defines the main conditionals.

```

3505 \cs_new_protected:Npn \__quark_new_conditional_n:Nnnnn
3506 { \__quark_new_conditional_aux_do:NNnnnn \use_i:nn }
3507 \cs_new_protected:Npn \__quark_new_conditional_N:Nnnnn
3508 { \__quark_new_conditional_aux_do:NNnnnn \use_ii:nn }

```

(End definition for __quark_new_conditional_n:Nnnnn and __quark_new_conditional_N:Nnnnn.)

__quark_new_conditional_aux_do:NNnnnn Similar to the previous macros, but branching conditionals only require one auxiliary, so we take a shortcut. In __quark_new_conditional_define:NNNNn, #4 is \use_i:nn to define the n-type function (which needs an auxiliary) and is \use_ii:nn to define the N-type function.

```

3509 \cs_new_protected:Npn \__quark_new_conditional_aux_do:NNnnnn #1 #2 #3 #4
3510 {
3511   \exp_args:Ncc \__quark_new_conditional_define:NNNNn
3512   { __ #4 _if_quark_ #3 :w } { q__ #4 _ #3 } #2 #1
3513 }
3514 \cs_new_protected:Npn \__quark_new_conditional_define:NNNNn #1 #2 #3 #4 #5
3515 {
3516   #4 { \cs_gset:Npn #1 ##1 #2 ##2 ? ##3 ?! { ##1 ##2 } } { }
3517   \exp_args:Nno \use:n { \prg_new_conditional:Npnn #3 ##1 {#5} }
3518   {
3519     #4 { \__quark_if_empty_if:o { #1 {} ##1 {} ?! #2 ??! } }
3520     { \if_meaning:w #2 ##1 }
3521     \prg_return_true: \else: \prg_return_false: \fi:
3522   }
3523 }

```

(End definition for __quark_new_conditional_aux_do:NNnnnn and __quark_new_conditional_define:NNNNn.)

__quark_module_name:N __quark_module_name:N takes a control sequence and returns its $\langle module \rangle$ name, determined as the first non-empty non-single-character word, separated by _ or :. These rules give the correct result for public functions \langle module \rangle_..., private functions __-\langle module \rangle_..., and variables such as \l_1\langle module \rangle_.... If no valid module is found the

result is an empty string. The approach is to first cut off everything after the (first) : if any is present, then repeatedly grab -delimited words until finding one of length at least 2 (we use low-level tests as `l3tl` is not fully available when `__kernel_quark_new_test:N` is first used. If no $\langle module \rangle$ is found (such as in `\::n`) we get the trailing marker `\use_none:n {}`, which expands to nothing.

```

3524 \cs_set:Npn \__quark_tmp:w #1#2
3525 {
3526   \cs_new:Npn \__quark_module_name:N ##1
3527   {
3528     \exp_last_unbraced:Nf \__quark_module_name:w
3529     { \cs_to_str:N ##1 } #1 \s__quark
3530   }
3531   \cs_new:Npn \__quark_module_name:w ##1 #1 ##2 \s__quark
3532   { \__quark_module_name_loop:w ##1 #2 \use_none:n { } #2 \s__quark }
3533   \cs_new:Npn \__quark_module_name_loop:w ##1 #2
3534   {
3535     \use_i_ii:nnn \if_meaning:w \prg_do_nothing:
3536     ##1 \prg_do_nothing: \prg_do_nothing:
3537     \exp_after:wN \__quark_module_name_loop:w
3538     \else:
3539     \__quark_module_name_end:w ##1
3540     \fi:
3541   }
3542   \cs_new:Npn \__quark_module_name_end:w
3543   ##1 \fi: ##2 \s__quark { \fi: ##1 }
3544 }
3545 \exp_after:wN \__quark_tmp:w \tl_to_str:n { : _ }

```

(End definition for `__quark_module_name:N` and others.)

`_quark_quark_conditional_name:N`
`_quark_quark_conditional_name:w`

`__quark_quark_conditional_name:N` determines the quark name that the quark conditional function `##1` queries, as the part of the function name between `_quark_if_` and the trailing `:`. Again we define it through `__quark_tmp:w`, which receives `:` as `#1` and `_quark_if_` as `#2`. The auxiliary `__quark_quark_conditional_name:w` returns the part between the first `_quark_if_` and the next `:`, and we apply this auxiliary to the function name followed by `:` (in case the function name is lacking a signature), and `_quark_if_:` so that `__quark_quark_conditional_name:N` returns an empty string if `_quark_if_` is not present.

```

3546 \cs_set:Npn \__quark_tmp:w #1 #2 \s__quark
3547 {
3548   \cs_new:Npn \__quark_quark_conditional_name:N ##1
3549   {
3550     \exp_last_unbraced:Nf \__quark_quark_conditional_name:w
3551     { \cs_to_str:N ##1 } #1 #2 #1 \s__quark
3552   }
3553   \cs_new:Npn \__quark_quark_conditional_name:w
3554   ##1 #2 ##2 #1 ##3 \s__quark {##2}
3555 }
3556 \exp_after:wN \__quark_tmp:w \tl_to_str:n { : _quark_if_ } \s__quark

```

(End definition for `__quark_quark_conditional_name:N` and `__quark_quark_conditional_name:w`.)

7.2 Scan marks

3557 <@@=scan>

\g__scan_marks_tl The list of all scan marks currently declared. No l3tl yet, so define this by hand.

3558 \cs_gset:Npn \g__scan_marks_tl { }

(End definition for \g__scan_marks_tl.)

\scan_new:N Check whether the variable is already a scan mark, then declare it to be equal to \scan_stop: globally.

```
3559 \cs_new_protected:Npn \scan_new:N #1
3560 {
3561   \tl_if_in:NnTF \g__scan_marks_tl { #1 }
3562   {
3563     \__kernel_msg_error:nxx { kernel } { scanmark-already-defined }
3564     { \token_to_str:N #1 }
3565   }
3566   {
3567     \tl_gput_right:Nn \g__scan_marks_tl {#1}
3568     \cs_new_eq:NN #1 \scan_stop:
3569   }
3570 }
```

(End definition for \scan_new:N. This function is documented on page 41.)

\s_stop We only declare one scan mark here, more can be defined by specific modules. Can't use \scan_new:N yet because l3tl isn't loaded, so define \s_stop by hand and add it to \g__scan_marks_tl. We also add \s__quark (declared earlier) to the pool here. Since it lives in a different namespace, a little l3docstrip cheating is necessary.

```
3571 \cs_new_eq:NN \s_stop \scan_stop:
3572 \cs_gset_nopar:Npx \g__scan_marks_tl
3573 {
3574   \exp_not:o \g__scan_marks_tl
3575   \s_stop
3576   <@@=quark>
3577   \s__quark
3578   <@@=scan>
3579 }
```

(End definition for \s_stop. This variable is documented on page 42.)

\use_none_delimit_by_s_stop:w Similar to \use_none_delimit_by_q_stop:w.

3580 \cs_new:Npn \use_none_delimit_by_s_stop:w #1 \s_stop { }

(End definition for \use_none_delimit_by_s_stop:w. This function is documented on page 42.)

3581 </package>

8 l3tl implementation

```
3582 \*package>
3583 \@@=tl>
```

A token list variable is a \TeX macro that holds tokens. By using the $\varepsilon\text{-TeX}$ primitive \unexpanded inside a \TeX \edef it is possible to store any tokens, including $\#$, in this way.

8.1 Functions

_kernel_tl_set:Nx These two are supplied to get better performance for macros which would otherwise use \tl_set:Nx or \tl_gset:Nx internally.

```
3584 \cs_new_eq:NN \_kernel\_tl\_set:Nx \cs_set_nopar:Npx
3585 \cs_new_eq:NN \_kernel\_tl\_gset:Nx \cs_gset_nopar:Npx
```

(End definition for _kernel_tl_set:Nx and $\text{_kernel_tl_gset:Nx}$.)

\tl_new:N Creating new token list variables is a case of checking for an existing definition and doing the definition.

```
\tl\_new:c
3586 \cs_new_protected:Npn \tl\_new:N #1
3587 {
3588   \_kernel\_chk_if_free_cs:N #1
3589   \cs_gset_eq:NN #1 \c\_empty\_tl
3590 }
3591 \cs_generate_variant:Nn \tl\_new:N { c }
```

(End definition for \tl_new:N . This function is documented on page 43.)

\tl_const:Nn Constants are also easy to generate. They use $\text{\cs_gset_nopar:Npx}$ instead of $\text{_kernel_tl_gset:Nx}$ so that the correct scope checking is applied if \l3debug is used.

```
\tl\_const:Nx
\tl\_const:cn
\tl\_const:cx
3592 \cs_new_protected:Npn \tl\_const:Nn #1#2
3593 {
3594   \_kernel\_chk_if_free_cs:N #1
3595   \cs_gset_nopar:Npx #1 { \_kernel\_exp\_not:w {#2} }
3596 }
3597 \cs_new_protected:Npn \tl\_const:Nx #1#2
3598 {
3599   \_kernel\_chk_if_free_cs:N #1
3600   \cs_gset_nopar:Npx #1 {#2}
3601 }
3602 \cs_generate_variant:Nn \tl\_const:Nn { c }
3603 \cs_generate_variant:Nn \tl\_const:Nx { c }
```

(End definition for \tl_const:Nn . This function is documented on page 43.)

\tl_clear:N Clearing a token list variable means setting it to an empty value. Error checking is sorted out by the parent function.

```
\tl\_clear:c
\tl\_gclear:N
\tl\_gclear:c
3604 \cs_new_protected:Npn \tl\_clear:N #1
3605 { \tl\_set_eq:NN #1 \c\_empty\_tl }
3606 \cs_new_protected:Npn \tl\_gclear:N #1
3607 { \tl\_gset_eq:NN #1 \c\_empty\_tl }
3608 \cs_generate_variant:Nn \tl\_clear:N { c }
3609 \cs_generate_variant:Nn \tl\_gclear:N { c }
```


(End definition for `\tl_clear:N` and `\tl_gclear:N`. These functions are documented on page 43.)

<code>\tl_clear_new:N</code>	Clearing a token list variable means setting it to an empty value. Error checking is sorted
<code>\tl_clear_new:c</code>	out by the parent function.

```

3610 \cs_new_protected:Npn \tl_clear_new:N #1
3611 { \tl_if_exist:NTF #1 { \tl_clear:N #1 } { \tl_new:N #1 } }
3612 \cs_new_protected:Npn \tl_gclear_new:N #1
3613 { \tl_if_exist:NTF #1 { \tl_gclear:N #1 } { \tl_new:N #1 } }
3614 \cs_generate_variant:Nn \tl_clear_new:N { c }
3615 \cs_generate_variant:Nn \tl_gclear_new:N { c }

```

(End definition for `\tl_clear_new:N` and `\tl_gclear_new:N`. These functions are documented on page 44.)

<code>\tl_set_eq:NN</code>	For setting token list variables equal to each other. To allow for patching, the arguments
<code>\tl_set_eq:Nc</code>	have to be explicit.

```

3616 \cs_new_protected:Npn \tl_set_eq:NN #1#2 { \cs_set_eq:NN #1 #2 }
3617 \cs_new_protected:Npn \tl_gset_eq:NN #1#2 { \cs_gset_eq:NN #1 #2 }
3618 \cs_generate_variant:Nn \tl_set_eq:NN { cN, Nc, cc }
3619 \cs_generate_variant:Nn \tl_gset_eq:NN { cN, Nc, cc }

```

(End definition for `\tl_set_eq:NN` and `\tl_gset_eq:NN`. These functions are documented on page 44.)

<code>\tl_concat:NNN</code>	Concatenating token lists is easy. When checking is turned on, all three arguments must
<code>\tl_concat:ccc</code>	be checked: a token list #2 or #3 equal to <code>\scan_stop:</code> would lead to problems later on.

```

3620 \cs_new_protected:Npn \tl_concat:NNN #1#2#3
3621 {
3622     \__kernel_tl_set:Nx #1
3623     {
3624         \__kernel_exp_not:w \exp_after:wN {#2}
3625         \__kernel_exp_not:w \exp_after:wN {#3}
3626     }
3627 }
3628 \cs_new_protected:Npn \tl_gconcat:NNN #1#2#3
3629 {
3630     \__kernel_tl_gset:Nx #1
3631     {
3632         \__kernel_exp_not:w \exp_after:wN {#2}
3633         \__kernel_exp_not:w \exp_after:wN {#3}
3634     }
3635 }
3636 \cs_generate_variant:Nn \tl_concat:NNN { ccc }
3637 \cs_generate_variant:Nn \tl_gconcat:NNN { ccc }

```

(End definition for `\tl_concat:NNN` and `\tl_gconcat:NNN`. These functions are documented on page 44.)

`\tl_if_exist_p:N` Copies of the `cs` functions defined in `l3basics`.

```
3638 \prg_new_eq_conditional:NNn \tl_if_exist:N \cs_if_exist:N { TF , T , F , p }
3639 \prg_new_eq_conditional:NNn \tl_if_exist:c \cs_if_exist:c { TF , T , F , p }
```

(End definition for `\tl_if_exist:NTF`. This function is documented on page 44.)

8.2 Constant token lists

\c_empty_tl Never full. We need to define that constant before using `\tl_new:N`.

```
3640 \tl_const:Nn \c_empty_tl { }
```

(End definition for `\c_empty_tl`. This variable is documented on page 58.)

\c_novalue_tl A special marker: as we don't have `\char_generate:nn` yet, has to be created the old-fashioned way.

```
3641 \group_begin:
3642 \tex_lccode:D 'A = '-
3643 \tex_lccode:D 'N = 'N
3644 \tex_lccode:D 'V = 'V
3645 \tex_lowercase:D
3646 {
3647   \group_end:
3648   \tl_const:Nn \c_novalue_tl { ANoValue- }
3649 }
```

(End definition for `\c_novalue_tl`. This variable is documented on page 58.)

\c_space_tl A space as a token list (as opposed to as a character).

```
3650 \tl_const:Nn \c_space_tl { ~ }
```

(End definition for `\c_space_tl`. This variable is documented on page 58.)

8.3 Adding to token list variables

\tl_set:Nn By using `\exp_not:n` token list variables can contain # tokens, which makes the token list registers provided by T_EX more or less redundant. The `\tl_set:No` version is done “by hand” as it is used quite a lot.

```
\tl_set:Nv 3651 \cs_new_protected:Npn \tl_set:Nn #1#2
\tl_set:Nf 3652 { \__kernel_tl_set:Nx #1 { \__kernel_exp_not:w {#2} } }
\tl_set:Nx 3653 \cs_new_protected:Npn \tl_set:No #1#2
\tl_set:cn 3654 { \__kernel_tl_set:Nx #1 { \__kernel_exp_not:w \exp_after:wN {#2} } }
\tl_set:cV 3655 \cs_new_protected:Npn \tl_set:Nx #1#2
\tl_set:cv 3656 { \__kernel_tl_set:Nx #1 {#2} }
\tl_set:co 3657 \cs_new_protected:Npn \tl_gset:Nn #1#2
\tl_set:cf 3658 { \__kernel_tl_gset:Nx #1 { \__kernel_exp_not:w {#2} } }
\tl_set:cx 3659 \cs_new_protected:Npn \tl_gset:No #1#2
3660 { \__kernel_tl_gset:Nx #1 { \__kernel_exp_not:w \exp_after:wN {#2} } }
\tl_gset:Nn 3661 \cs_new_protected:Npn \tl_gset:Nx #1#2
\tl_gset:Nv 3662 { \__kernel_tl_gset:Nx #1 {#2} }
\tl_gset:Nf 3663 \cs_generate_variant:Nn \tl_set:Nn { NV , Nv , Nf }
\tl_gset:No 3664 \cs_generate_variant:Nn \tl_set:Nn { c , cV , cv , cf }
\tl_gset:Nf 3665 \cs_generate_variant:Nn \tl_set:Nx { c }
\tl_gset:Nx 3666 \cs_generate_variant:Nn \tl_set:No { c }
\tl_gset:cn 3667 \cs_generate_variant:Nn \tl_gset:Nn { NV , Nv , Nf }
\tl_gset:cV 3668 \cs_generate_variant:Nn \tl_gset:Nn { c , cV , cv , cf }
\tl_gset:cv 3669 \cs_generate_variant:Nn \tl_gset:Nx { c }
\tl_gset:co 3670 \cs_generate_variant:Nn \tl_gset:No { c }
```

(End definition for `\tl_set:Nn` and `\tl_gset:Nn`. These functions are documented on page 44.)

\tl_gset:cx

```

\tl_put_left:Nn Adding to the left is done directly to gain a little performance.
\tl_put_left:NV
\tl_put_left:No
\tl_put_left:Nx
\tl_put_left:cn
\tl_put_left:cV
\tl_put_left:co
\tl_put_left:cx
\tl_gput_left:Nn
\tl_gput_left:NV
\tl_gput_left:No
\tl_gput_left:Nx
\tl_gput_left:cn
\tl_gput_left:cV
\tl_gput_left:co
\tl_gput_left:cx
3671 \cs_new_protected:Npn \tl_put_left:Nn #1#2
3672 {
3673   \__kernel_tl_set:Nx #1
3674   { \__kernel_exp_not:w {#2} \__kernel_exp_not:w \exp_after:wN {#1} }
3675 }
3676 \cs_new_protected:Npn \tl_put_left:NV #1#2
3677 {
3678   \__kernel_tl_set:Nx #1
3679   { \exp_not:V #2 \__kernel_exp_not:w \exp_after:wN {#1} }
3680 }
3681 \cs_new_protected:Npn \tl_put_left:No #1#2
3682 {
3683   \__kernel_tl_set:Nx #1
3684   {
3685     \__kernel_exp_not:w \exp_after:wN {#2}
3686     \__kernel_exp_not:w \exp_after:wN {#1}
3687   }
3688 }
3689 \cs_new_protected:Npn \tl_put_left:Nx #1#2
3690 { \__kernel_tl_set:Nx #1 { #2 \__kernel_exp_not:w \exp_after:wN {#1} } }
3691 \cs_new_protected:Npn \tl_gput_left:Nn #1#2
3692 {
3693   \__kernel_tl_gset:Nx #1
3694   { \__kernel_exp_not:w {#2} \__kernel_exp_not:w \exp_after:wN {#1} }
3695 }
3696 \cs_new_protected:Npn \tl_gput_left:NV #1#2
3697 {
3698   \__kernel_tl_gset:Nx #1
3699   { \exp_not:V #2 \__kernel_exp_not:w \exp_after:wN {#1} }
3700 }
3701 \cs_new_protected:Npn \tl_gput_left:No #1#2
3702 {
3703   \__kernel_tl_gset:Nx #1
3704   {
3705     \__kernel_exp_not:w \exp_after:wN {#2}
3706     \__kernel_exp_not:w \exp_after:wN {#1}
3707   }
3708 }
3709 \cs_new_protected:Npn \tl_gput_left:Nx #1#2
3710 { \__kernel_tl_gset:Nx #1 { #2 \__kernel_exp_not:w \exp_after:wN {#1} } }
3711 \cs_generate_variant:Nn \tl_put_left:Nn { c }
3712 \cs_generate_variant:Nn \tl_put_left:NV { c }
3713 \cs_generate_variant:Nn \tl_put_left:No { c }
3714 \cs_generate_variant:Nn \tl_put_left:Nx { c }
3715 \cs_generate_variant:Nn \tl_gput_left:Nn { c }
3716 \cs_generate_variant:Nn \tl_gput_left:NV { c }
3717 \cs_generate_variant:Nn \tl_gput_left:No { c }
3718 \cs_generate_variant:Nn \tl_gput_left:Nx { c }

```

(End definition for \tl_put_left:Nn and \tl_gput_left:Nn. These functions are documented on page 44.)

\tl_put_right:Nn The same on the right.

```

\tl_put_right:NV
\tl_put_right:No
\tl_put_right:Nx
\tl_put_right:cn
\tl_put_right:cV
\tl_put_right:co
\tl_put_right:cx
\tl_gput_right:Nn
\tl_gput_right:NV
\tl_gput_right:No
\tl_gput_right:Nx

```

```

3719 \cs_new_protected:Npn \tl_put_right:Nn #1#2
3720 { \__kernel_tl_set:Nx #1 { \__kernel_exp_not:w \exp_after:wN { #1 #2 } } }
3721 \cs_new_protected:Npn \tl_put_right:NV #1#2
3722 {
3723   \__kernel_tl_set:Nx #1
3724   { \__kernel_exp_not:w \exp_after:wN {#1} \exp_not:V #2 }
3725 }
3726 \cs_new_protected:Npn \tl_put_right:No #1#2
3727 {
3728   \__kernel_tl_set:Nx #1
3729   {
3730     \__kernel_exp_not:w \exp_after:wN {#1}
3731     \__kernel_exp_not:w \exp_after:wN {#2}
3732   }
3733 }
3734 \cs_new_protected:Npn \tl_put_right:Nx #1#2
3735 { \__kernel_tl_set:Nx #1 { \__kernel_exp_not:w \exp_after:wN {#1} #2 } }
3736 \cs_new_protected:Npn \tl_gput_right:Nn #1#2
3737 { \__kernel_tl_gset:Nx #1 { \__kernel_exp_not:w \exp_after:wN { #1 #2 } } }
3738 \cs_new_protected:Npn \tl_gput_right:NV #1#2
3739 {
3740   \__kernel_tl_gset:Nx #1
3741   { \__kernel_exp_not:w \exp_after:wN {#1} \exp_not:V #2 }
3742 }
3743 \cs_new_protected:Npn \tl_gput_right:No #1#2
3744 {
3745   \__kernel_tl_gset:Nx #1
3746   {
3747     \__kernel_exp_not:w \exp_after:wN {#1}
3748     \__kernel_exp_not:w \exp_after:wN {#2}
3749   }
3750 }
3751 \cs_new_protected:Npn \tl_gput_right:Nx #1#2
3752 { \__kernel_tl_gset:Nx #1 { \__kernel_exp_not:w \exp_after:wN {#1} #2 } }
3753 \cs_generate_variant:Nn \tl_put_right:Nn { c }
3754 \cs_generate_variant:Nn \tl_put_right:NV { c }
3755 \cs_generate_variant:Nn \tl_put_right:No { c }
3756 \cs_generate_variant:Nn \tl_put_right:Nx { c }
3757 \cs_generate_variant:Nn \tl_gput_right:Nn { c }
3758 \cs_generate_variant:Nn \tl_gput_right:NV { c }
3759 \cs_generate_variant:Nn \tl_gput_right:No { c }
3760 \cs_generate_variant:Nn \tl_gput_right:Nx { c }

```

(End definition for `\tl_put_right:Nn` and `\tl_gput_right:Nn`. These functions are documented on page 44.)

8.4 Internal quarks and quark-query functions

```

\q__tl_nil Internal quarks.
\q__tl_mark 3761 \quark_new:N \q__tl_nil
\q__tl_stop 3762 \quark_new:N \q__tl_mark
3763 \quark_new:N \q__tl_stop

```

(End definition for `\q__tl_nil`, `\q__tl_mark`, and `\q__tl_stop`.)

```
\q__tl_recursion_tail Internal recursion quarks.
\q__tl_recursion_stop 3764 \quark_new:N \q__tl_recursion_tail
3765 \quark_new:N \q__tl_recursion_stop

(End definition for \q__tl_recursion_tail and \q__tl_recursion_stop.)
```

```
\_tl_if_recursion_tail_break:nN Functions to query recursion quarks.
\_tl_if_recursion_tail_stop_p:n 3766 \__kernel_quark_new_test:N \_tl_if_recursion_tail_break:nN
\_tl_if_recursion_tail_stop:nTF 3767 \__kernel_quark_new_conditional:Nn \_tl_quark_if_nil:n { TF }

(End definition for \_tl_if_recursion_tail_break:nN and \_tl_if_recursion_tail_stop:nTF.)
```

8.5 Reassigning token list category codes

`\c__tl_rescan_marker_tl` The rescanning code needs a special token list containing the same character (chosen here to be a colon) with two different category codes: it cannot appear in the tokens being rescanned since all colons have the same category code.

```
3768 \tl_const:Nx \c__tl_rescan_marker_tl { : \token_to_str:N : }

(End definition for \c__tl_rescan_marker_tl.)
```

```
\tl_set_rescan:Nnn In a group, after some initial setup explained below and the user setup #3 (followed by
\tl_set_rescan:Nno \scan_stop: to be safe), there is a call to \__tl_set_rescan:NNN. This shared auxiliary
\tl_set_rescan:Nnx defined later distinguishes single-line and multi-line “files”. In the simplest case of multi-
\tl_set_rescan:cnn line files, it calls (with the same arguments) \__tl_set_rescan_multi:NNN, whose code
\tl_set_rescan:cno is included here to help understand the approach. This function rescans its argument #1,
\tl_set_rescan:cnx closes the group, and performs the assignment.
```

```
\tl_gset_rescan:Nnn One difficulty when rescanning is that \scantokens treats the argument as a file,
\tl_gset_rescan:Nno and without the correct settings a TEX error occurs:
```

```
\tl_gset_rescan:Nnx ! File ended while scanning definition of ...
\tl_gset_rescan:cnn
```

```
\tl_gset_rescan:cno A related minor issue is a warning due to opening a group before the \scantokens and
\tl_gset_rescan:cnx closing it inside that temporary file; we avoid that by setting \tracingnesting. The
\tl_rescan:nn standard solution to the “File ended” error is to grab the rescanned tokens as a delimited
\tl_rescan:nn argument of an auxiliary, here \__tl_rescan:NNw, that performs the assignment, then let
\tl_rescan:nn TEX “execute” the end of file marker. As usual in delimited arguments we use \prg_do_
\tl_rescan:nn nothing: to avoid stripping an outer set braces: this is removed by using o-expanding
\tl_rescan:nn assignments. The delimiter cannot appear within the rescanned token list because it
\tl_rescan:nn contains twice the same character, with different catcodes.
```

For `\tl_rescan:nn` we cannot simply call `__tl_set_rescan:NNnn \prg_do_nothing: \use:n` because that would leave the end-of-file marker *after* the result of rescanning. If that rescanned result is code that looks further in the input stream for arguments, it would break.

For multi-line files the only subtlety is that `\newlinechar` should be equal to `\endlinechar` because `\newlinechar` characters become new lines and then become `\endlinechar` characters when writing to an abstract file and reading back. This equality is ensured by setting `\newlinechar` equal to `\endlinechar`. Prior to this, `\endlinechar` is set to `-1` if it was `32` (in particular true after `\ExplSyntaxOn`) to avoid unreasonable line-breaks at every space for instance in error messages triggered by the user setup. Another side effect of reading back from the file is that spaces (catcode 10) are ignored

at the beginning of lines, and spaces and tabs (character code 32 and 9) are ignored at the end of lines.

The two `\if_false: ... \fi:` are there to prevent alignment tabs to cause a change of tabular cell while rescanning. We put the “opening” one after `\group_begin:` so that if one accidentally f-expands `\tl_set_rescan:Nnn` braces remain balanced. This is essential in e-type arguments when `\expanded` is not available.

```

3769 \cs_new_protected:Npn \tl_rescan:nn #1#2
3770 {
3771   \tl_set_rescan:Nnn \l__tl_internal_a_tl {#1} {#2}
3772   \exp_after:wN \__tl_rescan_aux:
3773   \l__tl_internal_a_tl
3774 }
3775 \exp_args:NNo \cs_new_protected:Npn \__tl_rescan_aux:
3776 { \tl_clear:N \l__tl_internal_a_tl }
3777 \cs_new_protected:Npn \tl_set_rescan:Nnn
3778 { \__tl_set_rescan:NNnn \tl_set:No }
3779 \cs_new_protected:Npn \tl_gset_rescan:Nnn
3780 { \__tl_set_rescan:NNnn \tl_gset:No }
3781 \cs_new_protected:Npn \__tl_set_rescan:NNnn #1#2#3#4
3782 {
3783   \group_begin:
3784   \if_false: { \fi:
3785     \int_set_eq:NN \tex_tracingnesting:D \c_zero_int
3786     \int_compare:nNnT \tex_endlinechar:D = { 32 }
3787     { \int_set:Nn \tex_endlinechar:D { -1 } }
3788     \int_set_eq:NN \tex_newlinechar:D \tex_endlinechar:D
3789     #3 \scan_stop:
3790     \exp_args:No \__tl_set_rescan:nNN { \tl_to_str:n {#4} } #1 #2
3791   \if_false: } \fi:
3792 }
3793 \cs_new_protected:Npn \__tl_set_rescan_multi:nNN #1#2#3
3794 {
3795   \tex_everyeof:D \exp_after:wN { \c__tl_rescan_marker_tl }
3796   \exp_after:wN \__tl_rescan:NNw
3797   \exp_after:wN #2
3798   \exp_after:wN #3
3799   \exp_after:wN \prg_do_nothing:
3800   \tex_scantokens:D {#1}
3801 }
3802 \exp_args:Nno \use:nn
3803 { \cs_new:Npn \__tl_rescan:NNw #1#2#3 } \c__tl_rescan_marker_tl
3804 {
3805   \group_end:
3806   #1 #2 {#3}
3807 }
3808 \cs_generate_variant:Nn \tl_set_rescan:Nnn { Nno , Nnx }
3809 \cs_generate_variant:Nn \tl_set_rescan:Nnn { c , cno , cnx }
3810 \cs_generate_variant:Nn \tl_gset_rescan:Nnn { Nno , Nnx }
3811 \cs_generate_variant:Nn \tl_gset_rescan:Nnn { c , cno }

```

(End definition for `\tl_set_rescan:Nnn` and others. These functions are documented on page 46.)

```

\__tl_set_rescan:nNN The function \__tl_set_rescan:nNN calls \__tl_set_rescan_multi:nNN or \__tl_-
\__tl_set_rescan_single:nNN set_rescan_single:nNN { ' } depending on whether its argument is a single-line
\__tl_set_rescan_single_aux:nnnNN
\__tl_set_rescan_single_aux:w

```

fragment of code/data or is made of multiple lines by testing for the presence of a `\newlinechar` character. If `\newlinechar` is out of range, the argument is assumed to be a single line.

For a single line, no `\endlinechar` should be added, so it is set to `-1`, and spaces should not be removed. Trailing spaces and tabs are a difficult matter, as `TeX` removes these at a very low level. The only way to preserve them is to rescan not the argument but the argument followed by a character with a reasonable category code. Here, `11` (letter) and `12` (other) are accepted, as these are convenient, suitable for delimiting an argument, and it is very unlikely that none of the ASCII characters are in one of these categories. To avoid selecting one particular character to put at the end, whose category code may have been modified, there is a loop through characters from `'` (ASCII 39) to `~` (ASCII 127). The choice of starting point was made because this is the start of a very long range of characters whose standard category is letter or other, thus minimizing the number of steps needed by the loop (most often just a single one). If no valid character is found (very rare), fall-back on `__tl_set_rescan_multi:nnn`.

Otherwise, once a valid character is found (let us use `'` in this explanation) run some code very similar to `__tl_set_rescan_multi:nnn` but with `'` added at both ends of the input. Of course, we need to define the auxiliary `__tl_set_rescan_single:NNww` on the fly to remove the additional `'` that is just before `::` (by which we mean `\c__tl_rescan_marker_tl`). Note that the argument must be delimited by `'` with the current catcode; this is done thanks to `\char_generate:nn`. Yet another issue is that the rescanned token list may contain a comment character, in which case the `'` we expected is not there. We fix this as follows: rather than just `::` we set `\everyeof` to `::{<code1>}` `'::<code2>}` `\s__tl_stop`. The auxiliary `__tl_set_rescan_single:NNww` runs the `o`-expanding assignment, expanding either `<code1>` or `<code2>` before its the main argument `#3`. In the typical case without comment character, `<code1>` is expanded, removing the leading `'`. In the rarer case with comment character, `<code2>` is expanded, calling `__tl_set_rescan_single_aux:w`, which removes the trailing `::{<code1>}` and the leading `'`.

```

3812 \cs_new_protected:Npn \__tl_set_rescan:nnn #1
3813 {
3814   \int_compare:nNnTF \tex_newlinechar:D < 0
3815   { \use_i:nn }
3816   {
3817     \exp_args:Nnf \tl_if_in:nnTF {#1}
3818     { \char_generate:nn { \tex_newlinechar:D } { 12 } }
3819   }
3820   { \__tl_set_rescan_multi:nnn }
3821   {
3822     \int_set:Nn \tex_endlinechar:D { -1 }
3823     \__tl_set_rescan_single:nnn { ' ' }
3824   }
3825   {#1}
3826 }
3827 \cs_new_protected:Npn \__tl_set_rescan_single:nnn #1
3828 {
3829   \int_compare:nNnTF
3830   { \char_value_catcode:n {#1} / 2 } = 6
3831   {
3832     \exp_args:Nof \__tl_set_rescan_single_aux:nnnn
3833     \c__tl_rescan_marker_tl

```

```

3834         { \char_generate:nn {#1} { \char_value_catcode:n {#1} } }
3835     }
3836     {
3837         \int_compare:nNnTF {#1} < { '\~ }
3838         {
3839             \exp_args:Nf \__tl_set_rescan_single:nnNN
3840             { \int_eval:n { #1 + 1 } }
3841         }
3842         { \__tl_set_rescan_multi:nnN }
3843     }
3844 }
3845 \cs_new_protected:Npn \__tl_set_rescan_single_aux:nnnNN #1#2#3#4#5
3846 {
3847     \tex_everyeof:D
3848     {
3849         #1 \use_none:n
3850         #2 #1 { \exp:w \__tl_set_rescan_single_aux:w }
3851         \s__tl_stop
3852     }
3853     \cs_set:Npn \__tl_rescan:NNw ##1##2##3 #2 #1 ##4 ##5 \s__tl_stop
3854     {
3855         \group_end:
3856         ##1 ##2 { ##4 ##3 }
3857     }
3858     \exp_after:wN \__tl_rescan:NNw
3859     \exp_after:wN #4
3860     \exp_after:wN #5
3861     \tex_scantokens:D { #2 #3 #2 }
3862 }
3863 \exp_args:Nno \use:nn
3864 { \cs_new:Npn \__tl_set_rescan_single_aux:w #1 }
3865 \c_tl_rescan_marker_tl #2
3866 { \use_i:nn \exp_end: #1 }

```

(End definition for `__tl_set_rescan:nnN` and others.)

8.6 Modifying token list variables

`\tl_replace_all:Nnn` All of the `replace` functions call `__tl_replace:NnNNNnn` with appropriate arguments. The first two arguments are explained later. The next controls whether the replacement function calls itself (`__tl_replace_next:w`) or stops (`__tl_replace_wrap:w`) after the first replacement. Next comes an x-type assignment function `\tl_set:Nx` or `\tl_gset:Nx` for local or global replacements. Finally, the three arguments $\langle tl\ var \rangle$ $\{ \langle pattern \rangle \}$ $\{ \langle replacement \rangle \}$ provided by the user. When describing the auxiliary functions below, we denote the contents of the $\langle tl\ var \rangle$ by $\langle token\ list \rangle$.

```

3867 \cs_new_protected:Npn \tl_replace_once:Nnn
3868 { \__tl_replace:NnNNNnn \q_tl_mark ? \__tl_replace_wrap:w \__kernel_tl_set:Nx }
3869 \cs_new_protected:Npn \tl_greplace_once:Nnn
3870 { \__tl_replace:NnNNNnn \q_tl_mark ? \__tl_replace_wrap:w \__kernel_tl_gset:Nx }
3871 \cs_new_protected:Npn \tl_replace_all:Nnn
3872 { \__tl_replace:NnNNNnn \q_tl_mark ? \__tl_replace_next:w \__kernel_tl_set:Nx }
3873 \cs_new_protected:Npn \tl_greplace_all:Nnn
3874 { \__tl_replace:NnNNNnn \q_tl_mark ? \__tl_replace_next:w \__kernel_tl_gset:Nx }

```



```

3875 \cs_generate_variant:Nn \tl_replace_once:Nnn { c }
3876 \cs_generate_variant:Nn \tl_greplace_once:Nnn { c }
3877 \cs_generate_variant:Nn \tl_replace_all:Nnn { c }
3878 \cs_generate_variant:Nn \tl_greplace_all:Nnn { c }

```

(End definition for `\tl_replace_all:Nnn` and others. These functions are documented on page 45.)

```

\__tl_replace:NnNNNnn
\__tl_replace_auxi:NnnNNNnn
\__tl_replace_auxii:nNNNnn
\__tl_replace_next:w
\__tl_replace_next_aux:w
\__tl_replace_wrap:w

```

To implement the actual replacement auxiliary `__tl_replace_auxii:nNNNnn` we need a $\langle \text{delimiter} \rangle$ with the following properties:

- all occurrences of the $\langle \text{pattern} \rangle$ #6 in “ $\langle \text{token list} \rangle \langle \text{delimiter} \rangle$ ” belong to the $\langle \text{token list} \rangle$ and have no overlap with the $\langle \text{delimiter} \rangle$,
- the first occurrence of the $\langle \text{delimiter} \rangle$ in “ $\langle \text{token list} \rangle \langle \text{delimiter} \rangle$ ” is the trailing $\langle \text{delimiter} \rangle$.

We first find the building blocks for the $\langle \text{delimiter} \rangle$, namely two tokens $\langle A \rangle$ and $\langle B \rangle$ such that $\langle A \rangle$ does not appear in #6 and #6 is not $\langle B \rangle$ (this condition is trivial if #6 has more than one token). Then we consider the delimiters “ $\langle A \rangle$ ” and “ $\langle A \rangle \langle A \rangle^n \langle B \rangle \langle A \rangle^n \langle B \rangle$ ”, for $n \geq 1$, where $\langle A \rangle^n$ denotes n copies of $\langle A \rangle$, and we choose as our $\langle \text{delimiter} \rangle$ the first one which is not in the $\langle \text{token list} \rangle$.

Every delimiter in the set obeys the first condition: #6 does not contain $\langle A \rangle$ hence cannot be overlapping with the $\langle \text{token list} \rangle$ and the $\langle \text{delimiter} \rangle$, and it cannot be within the $\langle \text{delimiter} \rangle$ since it would have to be in one of the two $\langle B \rangle$ hence be equal to this single token (or empty, but this is an error case filtered separately). Given the particular form of these delimiters, for which no prefix is also a suffix, the second condition is actually a consequence of the weaker condition that the $\langle \text{delimiter} \rangle$ we choose does not appear in the $\langle \text{token list} \rangle$. Additionally, the set of delimiters is such that a $\langle \text{token list} \rangle$ of n tokens can contain at most $O(n^{1/2})$ of them, hence we find a $\langle \text{delimiter} \rangle$ with at most $O(n^{1/2})$ tokens in a time at most $O(n^{3/2})$. Bear in mind that these upper bounds are reached only in very contrived scenarios: we include the case “ $\langle A \rangle$ ” in the list of delimiters to try, so that the $\langle \text{delimiter} \rangle$ is simply `\q__tl_mark` in the most common situation where neither the $\langle \text{token list} \rangle$ nor the $\langle \text{pattern} \rangle$ contains `\q__tl_mark`.

Let us now ahead, optimizing for this most common case. First, two special cases: an empty $\langle \text{pattern} \rangle$ #6 is an error, and if #1 is absent from both the $\langle \text{token list} \rangle$ #5 and the $\langle \text{pattern} \rangle$ #6 then we can use it as the $\langle \text{delimiter} \rangle$ through `__tl_replace_auxii:nNNNnn {#1}`. Otherwise, we end up calling `__tl_replace:NnNNNnn` repeatedly with the first two arguments `\q__tl_mark {?}`, `\? {??}`, `\?? {???}`, and so on, until #6 does not contain the control sequence #1, which we take as our $\langle A \rangle$. The argument #2 only serves to collect ? characters for #1. Note that the order of the tests means that the first two are done every time, which is wasteful (for instance, we repeatedly test for the emptiness of #6). However, this is rare enough not to matter. Finally, choose $\langle B \rangle$ to be `\q__tl_nil` or `\q__tl_stop` such that it is not equal to #6.

The `__tl_replace_auxi:NnnNNNnn` auxiliary receives $\{\langle A \rangle\}$ and $\{\langle A \rangle^n \langle B \rangle\}$ as its arguments, initially with $n = 1$. If “ $\langle A \rangle \langle A \rangle^n \langle B \rangle \langle A \rangle^n \langle B \rangle$ ” is in the $\langle \text{token list} \rangle$ then increase n and try again. Once it is not anymore in the $\langle \text{token list} \rangle$ we take it as our $\langle \text{delimiter} \rangle$ and pass this to the `auxii` auxiliary.

```

3879 \cs_new_protected:Npn \__tl_replace:NnNNNnn #1#2#3#4#5#6#7
3880 {
3881   \tl_if_empty:nTF {#6}
3882   {
3883     \__kernel_msg_error:nxx { kernel } { empty-search-pattern }

```

```

3884         { \tl_to_str:n {#7} }
3885     }
3886     {
3887         \tl_if_in:onTF { #5 #6 } {#1}
3888         {
3889             \tl_if_in:nnTF {#6} {#1}
3890             { \exp_args:Nc \__tl_replace:NnnNNnn {#2} {#2?} }
3891             {
3892                 \__tl_quark_if_nil:nTF {#6}
3893                 { \__tl_replace_auxi:NnnNNnn {#5} {#1} { #1 \q__tl_stop } }
3894                 { \__tl_replace_auxi:NnnNNnn {#5} {#1} { #1 \q__tl_nil } }
3895             }
3896         }
3897         { \__tl_replace_auxii:NnnNNnn {#1} }
3898         #3#4#5 {#6} {#7}
3899     }
3900 }
3901 \cs_new_protected:Npn \__tl_replace_auxi:NnnNNnn #1#2#3
3902 {
3903     \tl_if_in:NnTF #1 { #2 #3 #3 }
3904     { \__tl_replace_auxi:NnnNNnn #1 { #2 #3 } {#2} }
3905     { \__tl_replace_auxii:NnnNNnn { #2 #3 #3 } }
3906 }

```

The auxiliary `__tl_replace_auxii:NnnNNnn` receives the following arguments:

$\langle \textit{delimiter} \rangle$ $\langle \textit{function} \rangle$ $\langle \textit{assignment} \rangle$
 $\langle \textit{tl var} \rangle$ $\{ \langle \textit{pattern} \rangle \}$ $\{ \langle \textit{replacement} \rangle \}$

All of its work is done between `\group_align_safe_begin:` and `\group_align_safe_end:` to avoid issues in alignments. It does the actual replacement within `#3 #4 { ... }`, an x-expanding $\langle \textit{assignment} \rangle$ `#3` to the $\langle \textit{tl var} \rangle$ `#4`. The auxiliary `__tl_replace_next:w` is called, followed by the $\langle \textit{token list} \rangle$, some tokens including the $\langle \textit{delimiter} \rangle$ `#1`, followed by the $\langle \textit{pattern} \rangle$ `#5`. This auxiliary finds an argument delimited by `#5` (the presence of a trailing `#5` avoids runaway arguments) and calls `__tl_replace_wrap:w` to test whether this `#5` is found within the $\langle \textit{token list} \rangle$ or is the trailing one.

If on the one hand it is found within the $\langle \textit{token list} \rangle$, then `##1` cannot contain the $\langle \textit{delimiter} \rangle$ `#1` that we worked so hard to obtain, thus `__tl_replace_wrap:w` gets `##1` as its own argument `##1`, and protects it against the x-expanding assignment. It also finds `\exp_not:n` as `##2` and does nothing to it, thus letting through `\exp_not:n { \langle \textit{replacement} \rangle }` into the assignment. Note that `__tl_replace_next:w` and `__tl_replace_wrap:w` are always called followed by two empty brace groups. These are safe because no delimiter can match them. They prevent losing braces when grabbing delimited arguments, but require the use of `\exp_not:o` and `\use_none:nn`, rather than simply `\exp_not:n`. Afterwards, `__tl_replace_next:w` is called to repeat the replacement, or `__tl_replace_wrap:w` if we only want a single replacement. In this second case, `##1` is the $\langle \textit{remaining tokens} \rangle$ in the $\langle \textit{token list} \rangle$ and `##2` is some $\langle \textit{ending code} \rangle$ which ends the assignment and removes the trailing tokens `#5` using some `\if_false: { \fi: }` trickery because `#5` may contain any delimiter.

If on the other hand the argument `##1` of `__tl_replace_next:w` is delimited by the trailing $\langle \textit{pattern} \rangle$ `#5`, then `##1` is “`{ } { } { } \langle \textit{token list} \rangle \langle \textit{delimiter} \rangle { \langle \textit{ending code} \rangle }`”, hence `__tl_replace_wrap:w` finds “`{ } { } { } \langle \textit{token list} \rangle`” as `##1` and the $\langle \textit{ending code} \rangle$

as **##2**. It leaves the *token list* into the assignment and unbraces the *ending code* which removes what remains (essentially the *delimiter* and *replacement*).

```

3907 \cs_new_protected:Npn \__tl_replace_auxii:nNNNnn #1#2#3#4#5#6
3908 {
3909   \group_align_safe_begin:
3910   \cs_set:Npn \__tl_replace_wrap:w ##1 #1 ##2
3911     { \__kernel_exp_not:w \exp_after:wN { \use_none:nn ##1 } ##2 }
3912   \cs_set:Npx \__tl_replace_next:w ##1 #5
3913   {
3914     \exp_not:N \__tl_replace_wrap:w ##1
3915     \exp_not:n { #1 }
3916     \exp_not:n { \exp_not:n {#6} }
3917     \exp_not:n { #2 { } { } }
3918   }
3919   #3 #4
3920   {
3921     \exp_after:wN \__tl_replace_next_aux:w
3922     #4
3923     #1
3924     {
3925       \if_false: { \fi: }
3926       \exp_after:wN \use_none:n \exp_after:wN { \if_false: } \fi:
3927     }
3928     #5
3929   }
3930   \group_align_safe_end:
3931 }
3932 \cs_new:Npn \__tl_replace_next_aux:w { \__tl_replace_next:w { } { } }
3933 \cs_new_eq:NN \__tl_replace_wrap:w ?
3934 \cs_new_eq:NN \__tl_replace_next:w ?

```

(End definition for `__tl_replace:NnNNNnn` and others.)

<code>\tl_remove_once:Nn</code> <code>\tl_remove_once:cn</code> <code>\tl_gremove_once:Nn</code> <code>\tl_gremove_once:cn</code>	Removal is just a special case of replacement.	<pre> 3935 \cs_new_protected:Npn \tl_remove_once:Nn #1#2 3936 { \tl_replace_once:Nnn #1 {#2} { } } 3937 \cs_new_protected:Npn \tl_gremove_once:Nn #1#2 3938 { \tl_greplace_once:Nnn #1 {#2} { } } 3939 \cs_generate_variant:Nn \tl_remove_once:Nn { c } 3940 \cs_generate_variant:Nn \tl_gremove_once:Nn { c } </pre>
--	--	---

(End definition for `\tl_remove_once:Nn` and `\tl_gremove_once:Nn`. These functions are documented on page 45.)

<code>\tl_remove_all:Nn</code> <code>\tl_remove_all:cn</code> <code>\tl_gremove_all:Nn</code> <code>\tl_gremove_all:cn</code>	Removal is just a special case of replacement.	<pre> 3941 \cs_new_protected:Npn \tl_remove_all:Nn #1#2 3942 { \tl_replace_all:Nnn #1 {#2} { } } 3943 \cs_new_protected:Npn \tl_gremove_all:Nn #1#2 3944 { \tl_greplace_all:Nnn #1 {#2} { } } 3945 \cs_generate_variant:Nn \tl_remove_all:Nn { c } 3946 \cs_generate_variant:Nn \tl_gremove_all:Nn { c } </pre>
--	--	---

(End definition for `\tl_remove_all:Nn` and `\tl_gremove_all:Nn`. These functions are documented on page 45.)

8.7 Token list conditionals

`\tl_if_empty_p:N`
`\tl_if_empty_p:c`
`\tl_if_empty:NTF`
`\tl_if_empty:cTF`

These functions check whether the token list in the argument is empty and execute the proper code from their argument(s).

```

3947 \prg_new_conditional:Npnn \tl_if_empty:N #1 { p , T , F , TF }
3948 {
3949   \if_meaning:w #1 \c_empty_tl
3950   \prg_return_true:
3951   \else:
3952     \prg_return_false:
3953   \fi:
3954 }
3955 \prg_generate_conditional_variant:Nnn \tl_if_empty:N
3956 { c } { p , T , F , TF }

```

(End definition for `\tl_if_empty:NTF`. This function is documented on page 47.)

`\tl_if_empty_p:n`
`\tl_if_empty_p:V`
`\tl_if_empty:nTF`
`\tl_if_empty:VTF`

The `\if:w` triggers the expansion of `\tl_to_str:n` which converts the argument to a string: this is empty if and only if the argument is. Then `\if:w \scan_stop: ... \scan_stop:` is true if and only if the string ... is empty. It could be tempting to use `\if:w \scan_stop: #1 \scan_stop:` directly. But this fails on a token list expanding to anything starting with `\scan_stop:` leaving everything that follows in the input stream.

```

3957 \prg_new_conditional:Npnn \tl_if_empty:n #1 { p , TF , T , F }
3958 {
3959   \if:w \scan_stop: \tl_to_str:n {#1} \scan_stop:
3960   \prg_return_true:
3961   \else:
3962     \prg_return_false:
3963   \fi:
3964 }
3965 \prg_generate_conditional_variant:Nnn \tl_if_empty:n
3966 { V } { p , TF , T , F }

```

(End definition for `\tl_if_empty:nTF`. This function is documented on page 47.)

`\tl_if_empty_p:o`
`\tl_if_empty:oTF`
`__tl_if_empty_if:o`

The auxiliary function `__tl_if_empty_if:o` is for use in various token list conditionals which reduce to testing if a given token list is empty after applying a simple function to it. The test for emptiness is based on `\tl_if_empty:nTF`, but the expansion is hard-coded for efficiency, as this auxiliary function is used in several places. We don't put `\prg_return_true:` and so on in the definition of the auxiliary, because that would prevent an optimization applied to conditionals that end with this code. Also the `\@@_if_empty_if:o` is expanded once in `\tl_if_empty:oTF` for efficiency as well (and to reduce code doubling).

```

3967 \cs_new:Npn \__tl_if_empty_if:o #1
3968 {
3969   \if:w \scan_stop: \__kernel_tl_to_str:w \exp_after:wN {#1} \scan_stop:
3970 }
3971 \exp_args:Nno \use:n
3972 { \prg_new_conditional:Npnn \tl_if_empty:o #1 { p , TF , T , F } }
3973 {
3974   \__tl_if_empty_if:o {#1}
3975   \prg_return_true:

```

```

3976 \else:
3977 \prg_return_false:
3978 \fi:
3979 }

```

(End definition for `\tl_if_empty:nTF` and `__tl_if_empty_if:o`. This function is documented on page 47.)

```

\tl_if_blank_p:n TeX skips spaces when reading a non-delimited arguments. Thus, a  $\langle token list \rangle$  is blank
\tl_if_blank_p:V if and only if \use_none:n  $\langle token list \rangle$  ? is empty after one expansion. The auxiliary
\tl_if_blank_p:o \__tl_if_empty_if:o is a fast emptyness test, converting its argument to a string (after
\tl_if_blank:nTF one expansion) and using the test \if:w \scan_stop: ... \scan_stop:.
\tl_if_blank:VTF
\tl_if_blank:oTF
\__tl_if_blank_p:NNw
3980 \exp_args:Nno \use:n
3981 { \prg_new_conditional:Npnn \tl_if_blank:n #1 { p , T , F , TF } }
3982 {
3983 \__tl_if_empty_if:o { \use_none:n #1 ? }
3984 \prg_return_true:
3985 \else:
3986 \prg_return_false:
3987 \fi:
3988 }
3989 \prg_generate_conditional_variant:Nnn \tl_if_blank:n
3990 { e , V , o } { p , T , F , TF }

```

(End definition for `\tl_if_blank:nTF` and `__tl_if_blank_p:NNw`. This function is documented on page 46.)

```

\tl_if_eq_p:NN Returns \c_true_bool if and only if the two token list variables are equal.
\tl_if_eq_p:Nc
\tl_if_eq_p:cN
\tl_if_eq_p:cc
\tl_if_eq:NNTF
\tl_if_eq:NcTF
\tl_if_eq:cNTF
\tl_if_eq:ccTF
3991 \prg_new_conditional:Npnn \tl_if_eq:NN #1#2 { p , T , F , TF }
3992 {
3993 \if_meaning:w #1 #2
3994 \prg_return_true:
3995 \else:
3996 \prg_return_false:
3997 \fi:
3998 }
3999 \prg_generate_conditional_variant:Nnn \tl_if_eq:NN
4000 { Nc , c , cc } { p , TF , T , F }

```

(End definition for `\tl_if_eq:NNTF`. This function is documented on page 47.)

```

\l__tl_internal_a_tl Temporary storage.
\l__tl_internal_b_tl
4001 \tl_new:N \l__tl_internal_a_tl
4002 \tl_new:N \l__tl_internal_b_tl

```

(End definition for `\l__tl_internal_a_tl` and `\l__tl_internal_b_tl`.)

```

\tl_if_eq:NnTF A simple store and compare routine.
4003 \prg_new_protected_conditional:Npnn \tl_if_eq:Nn #1#2 { T , F , TF }
4004 {
4005 \group_begin:
4006 \tl_set:Nn \l__tl_internal_b_tl {#2}
4007 \exp_after:wN
4008 \group_end:
4009 \if_meaning:w #1 \l__tl_internal_b_tl

```

```

4010     \prg_return_true:
4011   \else:
4012     \prg_return_false:
4013   \fi:
4014 }
4015 \prg_generate_conditional_variant:Nnn \tl_if_eq:Nn { c } { TF , T , F }

```

(End definition for `\tl_if_eq:NnTF`. This function is documented on page 47.)

`\tl_if_eq:NnTF` A simple store and compare routine.

```

4016 \prg_new_protected_conditional:Npnn \tl_if_eq:Nn #1#2 { T , F , TF }
4017 {
4018   \group_begin:
4019   \tl_set:Nn \l__tl_internal_a_tl {#1}
4020   \tl_set:Nn \l__tl_internal_b_tl {#2}
4021   \exp_after:wN
4022   \group_end:
4023   \if_meaning:w \l__tl_internal_a_tl \l__tl_internal_b_tl
4024     \prg_return_true:
4025   \else:
4026     \prg_return_false:
4027   \fi:
4028 }

```

(End definition for `\tl_if_eq:NnTF`. This function is documented on page 47.)

`\tl_if_in:NnTF` See `\tl_if_in:nTF` for further comments. Here we simply expand the token list variable
`\tl_if_in:cnTF` and pass it to `\tl_if_in:nTF`.

```

4029 \cs_new_protected:Npn \tl_if_in:NnT { \exp_args:No \tl_if_in:nNT }
4030 \cs_new_protected:Npn \tl_if_in:NnF { \exp_args:No \tl_if_in:nNF }
4031 \cs_new_protected:Npn \tl_if_in:NnTF { \exp_args:No \tl_if_in:nTF }
4032 \prg_generate_conditional_variant:Nnn \tl_if_in:Nn
4033 { c } { T , F , TF }

```

(End definition for `\tl_if_in:NnTF`. This function is documented on page 47.)

`\tl_if_in:nTF` Once more, the test relies on the emptiness test for robustness. The function `__tl_tmp:w` removes tokens until the first occurrence of `#2`. If this does not appear in `#1`, then
`\tl_if_in:VnTF` the final `#2` is removed, leaving an empty token list. Otherwise some tokens remain, and
`\tl_if_in:onTF` the test is false. See `\tl_if_empty:nTF` for details on the emptiness test.
`\tl_if_in:noTF`

Treating correctly cases like `\tl_if_in:nTF {a state}{states}`, where `#1#2` contains `#2` before the end, requires special care. To cater for this case, we insert `{ }{ }` between the two token lists. This marker may not appear in `#2` because of \TeX limitations on what can delimit a parameter, hence we are safe. Using two brace groups makes the test work also for empty arguments. The `\if_false:` constructions are a faster way to do `\group_align_safe_begin:` and `\group_align_safe_end:`. The `\scan_stop:` ensures that f-expanding `\tl_if_in:n` does not lead to unbalanced braces.

```

4034 \prg_new_protected_conditional:Npnn \tl_if_in:n #1#2 { T , F , TF }
4035 {
4036   \scan_stop:
4037   \if_false: { \fi:
4038     \cs_set:Npn \__tl_tmp:w ##1 #2 { }
4039     \tl_if_empty:oTF { \__tl_tmp:w #1 {} {} #2 }
4040     { \prg_return_false: } { \prg_return_true: }

```

```

4041     \if_false: } \fi:
4042   }
4043 \prg_generate_conditional_variant:Nnn \tl_if_in:nn
4044   { V , o , no } { T , F , TF }

```

(End definition for `\tl_if_in:nnTF`. This function is documented on page 47.)

`\tl_if_novalue_p:n` Tests for `-NoValue-`: this is similar to `\tl_if_in:nn` but set up to be expandable and
`\tl_if_novalue:nTF` to check the value exactly. The question mark prevents the auxiliary from losing braces.
`__tl_if_novalue:w`

```

4045 \cs_set_protected:Npn \__tl_tmp:w #1
4046 {
4047   \prg_new_conditional:Npnn \tl_if_novalue:n ##1
4048     { p , T , F , TF }
4049     {
4050       \str_if_eq:onTF
4051         { \__tl_if_novalue:w ? ##1 { } #1 }
4052         { ? { } #1 }
4053         { \prg_return_true: }
4054         { \prg_return_false: }
4055       }
4056     \cs_new:Npn \__tl_if_novalue:w ##1 #1 {##1}
4057   }
4058 \exp_args:No \__tl_tmp:w { \c_novalue_tl }

```

(End definition for `\tl_if_novalue:nTF` and `__tl_if_novalue:w`. This function is documented on page 48.)

`\tl_if_single_p:N` Expand the token list and feed it to `\tl_if_single:n`.

```

\tl_if_single:NTF
4059 \cs_new:Npn \tl_if_single_p:N { \exp_args:No \tl_if_single:n }
4060 \cs_new:Npn \tl_if_single:NT { \exp_args:No \tl_if_single:nT }
4061 \cs_new:Npn \tl_if_single:NF { \exp_args:No \tl_if_single:nF }
4062 \cs_new:Npn \tl_if_single:NTF { \exp_args:No \tl_if_single:nTF }

```

(End definition for `\tl_if_single:NTF`. This function is documented on page 48.)

`\tl_if_single_p:n` This test is similar to `\tl_if_empty:nTF`. Expanding `\use_none:nn #1 ??` once yields
`\tl_if_single:nTF` an empty result if #1 is blank, a single ? if #1 has a single item, and otherwise yields
`__tl_if_single:nnw` some tokens ending with ??. Then, `__kernel_tl_to_str:w` makes sure there are no
odd category codes. An earlier version would compare the result to a single ? using string
comparison, but the Lua call is slow in LuaTeX. Instead, `__tl_if_single:nnw` picks
the second token in front of it. If #1 is empty, this token is the trailing ? and the `\if:w`
test yields `false`. If #1 has a single item, the token is `\scan_stop:` and the `\if:w` test
yields `true`. Otherwise, it is one of the characters resulting from `\tl_to_str:n`, and the
`\if:w` test yields `false`. Note that `\if:w` and `__kernel_tl_to_str:w` are primitives
that take care of expansion.

```

4063 \prg_new_conditional:Npnn \tl_if_single:n #1 { p , T , F , TF }
4064 {
4065   \if:w \scan_stop: \exp_after:wN \__tl_if_single:nnw
4066     \__kernel_tl_to_str:w
4067     \exp_after:wN { \use_none:nn #1 ?? } \scan_stop: ? \s__tl_stop
4068   \prg_return_true:
4069   \else:
4070     \prg_return_false:
4071   \fi:

```

```

4072 }
4073 \cs_new:Npn \__tl_if_single:nnw #1#2#3 \s__tl_stop {#2}

```

(End definition for `\tl_if_single:nTF` and `__tl_if_single:nnw`. This function is documented on page 48.)

`\tl_if_single_token_p:n` There are four cases: empty token list, token list starting with a normal token, with a brace group, or with a space token. If the token list starts with a normal token, remove it and check for emptiness. For the next case, an empty token list is not a single token. Finally, we have a non-empty token list starting with a space or a brace group. Applying f-expansion yields an empty result if and only if the token list is a single space.

`\tl_if_single_token:nTF`

```

4074 \prg_new_conditional:Npnn \tl_if_single_token:n #1 { p , T , F , TF }
4075 {
4076   \tl_if_head_is_N_type:nTF {#1}
4077   { \__tl_if_empty_if:o { \use_none:n #1 } }
4078   {
4079     \tl_if_empty:nTF {#1}
4080     { \if_false: }
4081     { \__tl_if_empty_if:o { \exp:w \exp_end_continue_f:w #1 } }
4082   }
4083   \prg_return_true:
4084 \else:
4085   \prg_return_false:
4086 \fi:
4087 }

```

(End definition for `\tl_if_single_token:nTF`. This function is documented on page 48.)

`\tl_case:Nn`

`\tl_case:cn`

`\tl_case:NnTF`

`\tl_case:cnTF`

`__tl_case:nnTF`

`__tl_case:Nw`

`__tl_case_end:nw`

The aim here is to allow the case statement to be evaluated using a known number of expansion steps (two), and without needing to use an explicit “end of recursion” marker. That is achieved by using the test input as the final case, as this is always true. The trick is then to tidy up the output such that the appropriate case code plus either the true or false branch code is inserted.

```

4088 \cs_new:Npn \tl_case:Nn #1#2
4089 {
4090   \exp:w
4091   \__tl_case:NnTF #1 {#2} { } { }
4092 }
4093 \cs_new:Npn \tl_case:NnT #1#2#3
4094 {
4095   \exp:w
4096   \__tl_case:NnTF #1 {#2} {#3} { }
4097 }
4098 \cs_new:Npn \tl_case:NnF #1#2#3
4099 {
4100   \exp:w
4101   \__tl_case:NnTF #1 {#2} { } {#3}
4102 }
4103 \cs_new:Npn \tl_case:NnTF #1#2
4104 {
4105   \exp:w
4106   \__tl_case:NnTF #1 {#2}
4107 }
4108 \cs_new:Npn \__tl_case:NnTF #1#2#3#4

```



```

4109 { \_tl\_case:Nw #1 #2 #1 { } \s\_tl\_mark {#3} \s\_tl\_mark {#4} \s\_tl\_stop }
4110 \cs\_new:Npn \_tl\_case:Nw #1#2#3
4111 {
4112   \tl\_if\_eq:NNTF #1 #2
4113   { \_tl\_case\_end:nw {#3} }
4114   { \_tl\_case:Nw #1 }
4115 }
4116 \cs\_generate\_variant:Nn \tl\_case:Nn { c }
4117 \prg\_generate\_conditional\_variant:Nnn \tl\_case:Nn
4118 { c } { T , F , TF }

```

To tidy up the recursion, there are two outcomes. If there was a hit to one of the cases searched for, then #1 is the code to insert, #2 is the *next* case to check on and #3 is all of the rest of the cases code. That means that #4 is the `true` branch code, and #5 tidies up the spare `\s_tl_mark` and the `false` branch. On the other hand, if none of the cases matched then we arrive here using the “termination” case of comparing the search with itself. That means that #1 is empty, #2 is the first `\s_tl_mark` and so #4 is the `false` code (the `true` code is mopped up by #3).

```

4119 \cs\_new:Npn \_tl\_case\_end:nw #1#2#3 \s\_tl\_mark #4#5 \s\_tl\_stop
4120 { \exp\_end: #1 #4 }

```

(End definition for `\tl_case:NnTF` and others. This function is documented on page 48.)

8.8 Mapping to token lists

`\tl_map_function:nN` Expandable loop macro for token lists. These have the advantage of not needing to test if the argument is empty, because if it is, the stop marker is read immediately and the loop terminated.

```

\__tl\_map\_function:Nn
4121 \cs\_new:Npn \tl\_map\_function:nN #1#2
4122 {
4123   \_tl\_map\_function:Nn #2 #1
4124   \q\_tl\_recursion\_tail
4125   \prg\_break\_point:Nn \tl\_map\_break: { }
4126 }
4127 \cs\_new:Npn \tl\_map\_function:NN
4128 { \exp\_args:No \tl\_map\_function:nN }
4129 \cs\_new:Npn \_tl\_map\_function:Nn #1#2
4130 {
4131   \_tl\_if\_recursion\_tail\_break:nN {#2} \tl\_map\_break:
4132   #1 {#2} \_tl\_map\_function:Nn #1
4133 }
4134 \cs\_generate\_variant:Nn \tl\_map\_function:NN { c }

```

(End definition for `\tl_map_function:nN`, `\tl_map_function:NN`, and `_tl_map_function:Nn`. These functions are documented on page 49.)

`\tl_map_inline:nn` The inline functions are straight forward by now. We use a little trick with the counter `\g_kernel_prg_map_int` to make them nestable. We can also make use of `_tl_map_function:Nn` from before.

```

4135 \cs\_new\_protected:Npn \tl\_map\_inline:nn #1#2
4136 {
4137   \int\_gincr:N \g\_kernel\_prg\_map\_int
4138   \cs\_gset\_protected:cpn
4139   { \_tl\_map\_ \int\_use:N \g\_kernel\_prg\_map\_int :w } ##1 {#2}

```

```

4140 \exp_args:Nc \__tl_map_function:Nn
4141 { __tl_map_ \int_use:N \g__kernel_prg_map_int :w }
4142 #1 \q__tl_recursion_tail
4143 \prg_break_point:Nn \tl_map_break:
4144 { \int_gdecr:N \g__kernel_prg_map_int }
4145 }
4146 \cs_new_protected:Npn \tl_map_inline:Nn
4147 { \exp_args:No \tl_map_inline:nn }
4148 \cs_generate_variant:Nn \tl_map_inline:Nn { c }

```

(End definition for `\tl_map_inline:nn` and `\tl_map_inline:Nn`. These functions are documented on page 49.)

`\tl_map_tokens:nn`
`\tl_map_tokens:Nn`
`\tl_map_tokens:cn`
`__tl_map_tokens:nn`

Much like the function mapping.

```

4149 \cs_new:Npn \tl_map_tokens:nn #1#2
4150 {
4151   \__tl_map_tokens:nn {#2} #1
4152   \q__tl_recursion_tail
4153   \prg_break_point:Nn \tl_map_break: { }
4154 }
4155 \cs_new:Npn \tl_map_tokens:Nn
4156 { \exp_args:No \tl_map_tokens:nn }
4157 \cs_generate_variant:Nn \tl_map_tokens:Nn { c }
4158 \cs_new:Npn \__tl_map_tokens:nn #1#2
4159 {
4160   \__tl_if_recursion_tail_break:nN {#2} \tl_map_break:
4161   \use:n {#1} {#2}
4162   \__tl_map_tokens:nn {#1}
4163 }

```

(End definition for `\tl_map_tokens:nn`, `\tl_map_tokens:Nn`, and `__tl_map_tokens:nn`. These functions are documented on page 49.)

`\tl_map_variable:nNn`
`\tl_map_variable:NNn`
`\tl_map_variable:cn`
`__tl_map_variable:Nnn`

`\tl_map_variable:nNn` *<token list>* *<tl var>* *<action>* assigns *<tl var>* to each element and executes *<action>*. The assignment to *<tl var>* is done after the quark test so that this variable does not get set to a quark.

```

4164 \cs_new_protected:Npn \tl_map_variable:nNn #1#2#3
4165 {
4166   \__tl_map_variable:Nnn #2 {#3} #1
4167   \q__tl_recursion_tail
4168   \prg_break_point:Nn \tl_map_break: { }
4169 }
4170 \cs_new_protected:Npn \tl_map_variable:NNn
4171 { \exp_args:No \tl_map_variable:nNn }
4172 \cs_new_protected:Npn \__tl_map_variable:Nnn #1#2#3
4173 {
4174   \__tl_if_recursion_tail_break:nN {#3} \tl_map_break:
4175   \tl_set:Nn #1 {#3}
4176   \use:n {#2}
4177   \__tl_map_variable:Nnn #1 {#2}
4178 }
4179 \cs_generate_variant:Nn \tl_map_variable:NNn { c }

```

(End definition for `\tl_map_variable:nNn`, `\tl_map_variable:NNn`, and `__tl_map_variable:Nnn`. These functions are documented on page 49.)

\tl_map_break: The break statements use the general \prg_map_break:Nn.
\tl_map_break:n

```
4180 \cs_new:Npn \tl_map_break:
4181 { \prg_map_break:Nn \tl_map_break: { } }
4182 \cs_new:Npn \tl_map_break:n
4183 { \prg_map_break:Nn \tl_map_break: }
```

(End definition for \tl_map_break: and \tl_map_break:n. These functions are documented on page 50.)

8.9 Using token lists

\tl_to_str:n Another name for a primitive: defined in l3basics.
\tl_to_str:V

```
4184 \cs_generate_variant:Nn \tl_to_str:n { V }
```

(End definition for \tl_to_str:n. This function is documented on page 51.)

\tl_to_str:N These functions return the replacement text of a token list as a string.

\tl_to_str:c

```
4185 \cs_new:Npn \tl_to_str:N #1 { \__kernel_tl_to_str:w \exp_after:wN {#1} }
4186 \cs_generate_variant:Nn \tl_to_str:N { c }
```

(End definition for \tl_to_str:N. This function is documented on page 51.)

\tl_use:N Token lists which are simply not defined give a clear T_EX error here. No such luck for
\tl_use:c ones equal to \scan_stop: so instead a test is made and if there is an issue an error is forced.

```
4187 \cs_new:Npn \tl_use:N #1
4188 {
4189   \tl_if_exist:NTF #1 {#1}
4190   {
4191     \__kernel_msg_expandable_error:nnn
4192     { kernel } { bad-variable } {#1}
4193   }
4194 }
4195 \cs_generate_variant:Nn \tl_use:N { c }
```

(End definition for \tl_use:N. This function is documented on page 51.)

8.10 Working with the contents of token lists

\tl_count:n Count number of elements within a token list or token list variable. Brace groups within
\tl_count:V the list are read as a single element. Spaces are ignored. __tl_count:n grabs the
\tl_count:o element and replaces it by +1. The 0 ensures that it works on an empty list.

\tl_count:N

```
4196 \cs_new:Npn \tl_count:n #1
4197 {
4198   \int_eval:n
4199   { 0 \tl_map_function:nN {#1} \__tl_count:n }
4200 }
4201 \cs_new:Npn \tl_count:N #1
4202 {
4203   \int_eval:n
4204   { 0 \tl_map_function:NN #1 \__tl_count:n }
4205 }
4206 \cs_new:Npn \__tl_count:n #1 { + 1 }
4207 \cs_generate_variant:Nn \tl_count:n { V , o }
4208 \cs_generate_variant:Nn \tl_count:N { c }
```

(End definition for `\tl_count:n`, `\tl_count:N`, and `__tl_count:n`. These functions are documented on page 51.)

`\tl_count_tokens:n` The token count is computed through an `\int_eval:n` construction. Each `1+` is output to the *left*, into the integer expression, and the sum is ended by the `\exp_end:` inserted by `__tl_act_end:wn` (which is technically implemented as `\c_zero_int`). Somewhat a hack!

```

4209 \cs_new:Npn \tl_count_tokens:n #1
4210 {
4211   \int_eval:n
4212   {
4213     \__tl_act:NNNn
4214     \__tl_act_count_normal:N
4215     \__tl_act_count_group:n
4216     \__tl_act_count_space:
4217     {#1}
4218   }
4219 }
4220 \cs_new:Npn \__tl_act_count_normal:N #1 { 1 + }
4221 \cs_new:Npn \__tl_act_count_space: { 1 + }
4222 \cs_new:Npn \__tl_act_count_group:n #1 { 2 + \tl_count_tokens:n {#1} + }

```

(End definition for `\tl_count_tokens:n` and others. This function is documented on page 52.)

`\tl_reverse_items:n` Reversal of a token list is done by taking one item at a time and putting it after `\s__tl_stop`.

```

4223 \cs_new:Npn \tl_reverse_items:n #1
4224 {
4225   \__tl_reverse_items:nwNwn #1 ?
4226   \s__tl_mark \__tl_reverse_items:nwNwn
4227   \s__tl_mark \__tl_reverse_items:wn
4228   \s__tl_stop { }
4229 }
4230 \cs_new:Npn \__tl_reverse_items:nwNwn #1 #2 \s__tl_mark #3 #4 \s__tl_stop #5
4231 {
4232   #3 #2
4233   \s__tl_mark \__tl_reverse_items:nwNwn
4234   \s__tl_mark \__tl_reverse_items:wn
4235   \s__tl_stop { {#1} #5 }
4236 }
4237 \cs_new:Npn \__tl_reverse_items:wn #1 \s__tl_stop #2
4238 { \__kernel_exp_not:w \exp_after:wN { \use_none:nn #2 } }

```

(End definition for `\tl_reverse_items:n`, `__tl_reverse_items:nwNwn`, and `__tl_reverse_items:wn`. This function is documented on page 52.)

`\tl_trim_spaces:n` Trimming spaces from around the input is deferred to an internal function whose first argument is the token list to trim, augmented by an initial `__tl_trim_mark:`, and whose second argument is a *<continuation>*, which receives as a braced argument `__tl_trim_mark: <trimmed token list>`. The control sequence `__tl_trim_mark:` expands to nothing in a single expansion. In the case at hand, we take `__kernel_exp_not:w \exp_after:wN` as our continuation, so that space trimming behaves correctly within an x-type expansion.

```

\__tl_trim_spaces:n
\__tl_trim_spaces:o
\__tl_trim_spaces_apply:n
\__tl_trim_spaces_apply:o
\__tl_trim_spaces:N
\__tl_trim_spaces:c
\__tl_gtrim_spaces:N
\__tl_gtrim_spaces:c
4239 \cs_new:Npn \tl_trim_spaces:n #1

```

```

4240 {
4241   \_tl_trim_spaces:nn
4242   { \_tl_trim_mark: #1 }
4243   { \_kernel_exp_not:w \exp_after:wN }
4244 }
4245 \cs_generate_variant:Nn \tl_trim_spaces:n { o }
4246 \cs_new:Npn \tl_trim_spaces_apply:nN #1#2
4247 { \_tl_trim_spaces:nn { \_tl_trim_mark: #1 } { \exp_args:No #2 } }
4248 \cs_generate_variant:Nn \tl_trim_spaces_apply:nN { o }
4249 \cs_new_protected:Npn \tl_trim_spaces:N #1
4250 { \_kernel_tl_set:Nx #1 { \exp_args:No \tl_trim_spaces:n {#1} } }
4251 \cs_new_protected:Npn \tl_gtrim_spaces:N #1
4252 { \_kernel_tl_gset:Nx #1 { \exp_args:No \tl_trim_spaces:n {#1} } }
4253 \cs_generate_variant:Nn \tl_trim_spaces:N { c }
4254 \cs_generate_variant:Nn \tl_gtrim_spaces:N { c }

```

Trimming spaces from around the input is done using delimited arguments and quarks, and to get spaces at odd places in the definitions, we nest those in `_tl_tmp:w`, which then receives a single space as its argument: `#1` is `␣`. Removing leading spaces is done with `_tl_trim_spaces_auxi:w`, which loops until `_tl_trim_mark:␣` matches the end of the token list: then `##1` is the token list and `##3` is `_tl_trim_spaces_auxii:w`. This hands the relevant tokens to the loop `_tl_trim_spaces_auxiii:w`, responsible for trimming trailing spaces. The end is reached when `␣ \s__tl_nil` matches the one present in the definition of `\tl_trim_spaces:n`. Then `_tl_trim_spaces_auxiv:w` puts the token list into a group, with a lingering `_tl_trim_mark:` at the start (which will expand to nothing in one step of expansion), and feeds this to the *continuation*.

```

4255 \cs_set_protected:Npn \_tl_tmp:w #1
4256 {
4257   \cs_new:Npn \_tl_trim_spaces:nn ##1
4258   {
4259     \_tl_trim_spaces_auxi:w
4260     ##1
4261     \s__tl_nil
4262     \_tl_trim_mark: #1 { }
4263     \_tl_trim_mark: \_tl_trim_spaces_auxii:w
4264     \_tl_trim_spaces_auxiii:w
4265     #1 \s__tl_nil
4266     \_tl_trim_spaces_auxiv:w
4267     \s__tl_stop
4268   }
4269   \cs_new:Npn
4270   \_tl_trim_spaces_auxi:w ##1 \_tl_trim_mark: #1 ##2 \_tl_trim_mark: ##3
4271   {
4272     ##3
4273     \_tl_trim_spaces_auxi:w
4274     \_tl_trim_mark:
4275     ##2
4276     \_tl_trim_mark: #1 {##1}
4277   }
4278   \cs_new:Npn \_tl_trim_spaces_auxii:w
4279   \_tl_trim_spaces_auxi:w \_tl_trim_mark: \_tl_trim_mark: ##1
4280   {
4281     \_tl_trim_spaces_auxiii:w

```

```

4282     ##1
4283   }
4284   \cs_new:Npn \__tl_trim_spaces_auxiii:w ##1 #1 \s__tl_nil ##2
4285   {
4286     ##2
4287     ##1 \s__tl_nil
4288     \__tl_trim_spaces_auxiii:w
4289   }
4290   \cs_new:Npn \__tl_trim_spaces_auxiv:w ##1 \s__tl_nil ##2 \s__tl_stop ##3
4291   { ##3 { ##1 } }
4292   \cs_new:Npn \__tl_trim_mark: {}
4293 }
4294 \__tl_tmp:w { ~ }

```

(End definition for `\tl_trim_spaces:n` and others. These functions are documented on page 52.)

`\tl_sort:Nn` Implemented in `l3sort`.

`\tl_sort:cn`

`\tl_gsort:Nn` (End definition for `\tl_sort:Nn`, `\tl_gsort:Nn`, and `\tl_sort:nN`. These functions are documented on page 53.)

`\tl_gsort:cn`

`\tl_sort:nN`

8.11 The first token from a token list

`\tl_head:N`

`\tl_head:n`

`\tl_head:V`

`\tl_head:v`

`\tl_head:f`

`__tl_head_auxi:nw`

`__tl_head_auxii:n`

`\tl_head:w`

`__tl_tl_head:w`

`\tl_tail:N`

`\tl_tail:n`

`\tl_tail:V`

`\tl_tail:v`

`\tl_tail:f`

Finding the head of a token list expandably always strips braces, which is fine as this is consistent with for example mapping to a list. The empty brace groups in `\tl_head:n` ensure that a blank argument gives an empty result. The result is returned within the `\unexpanded` primitive. The approach here is to use `\if_false:` to allow us to use `}` as the closing delimiter: this is the only safe choice, as any other token would not be able to parse it's own code. If the `\expanded` primitive is available it is used to get a fast and safe code variant in which we don't have to ensure that the left-most token is an internal to not break in an f-type expansion. If `\expanded` isn't available, using a marker, we can see if what we are grabbing is exactly the marker, or there is anything else to deal with. If there is, there is a loop. If not, tidy up and leave the item in the output stream. More detail in <http://tex.stackexchange.com/a/70168>.

```

4295 \cs_if_exist:NTF \tex_expanded:D
4296 {
4297   \cs_new:Npn \tl_head:n #1
4298   {
4299     \__kernel_exp_not:w \tex_expanded:D
4300     { { \if_false: { \fi: \__tl_head_aux:n #1 { } } } }
4301   }
4302   \cs_new:Npn \__tl_head_aux:n #1
4303   {
4304     \__kernel_exp_not:w {#1}
4305     \exp_after:wN \use_none:n \exp_after:wN { \if_false: } \fi:
4306   }
4307 }
4308 {
4309   \cs_new:Npn \tl_head:n #1
4310   {
4311     \__kernel_exp_not:w
4312     \if_false: { \fi: \__tl_head_auxi:nw #1 { } \s__tl_stop }
4313   }
4314   \cs_new:Npn \__tl_head_auxi:nw #1#2 \s__tl_stop

```

```

4315     {
4316       \exp_after:wN \_tl_head_auxii:n \exp_after:wN {
4317         \if_false: } \fi: {#1}
4318     }
4319     \exp_args:Nno \use:n
4320     { \cs_new:Npn \_tl_head_auxii:n #1 }
4321     {
4322       \_tl_if_empty_if:o { \use_none:n #1 }
4323       \exp_after:wN \use_ii:nnn
4324       \fi:
4325       \use_ii:nn
4326       {#1}
4327       { \if_false: { \fi: \_tl_head_auxi:nw #1 } }
4328     }
4329   }
4330   \cs_generate_variant:Nn \tl_head:n { V , v , f }
4331   \cs_new:Npn \tl_head:w #1#2 \q_stop {#1}
4332   \cs_new:Npn \_tl_tl_head:w #1#2 \s__tl_stop {#1}
4333   \cs_new:Npn \tl_head:N { \exp_args:No \tl_head:n }

```

To correctly leave the tail of a token list, it's important *not* to absorb any of the tail part as an argument. For example, the simple definition

```

\cs_new:Npn \tl_tail:n #1 { \tl_tail:w #1 \q_stop }
\cs_new:Npn \tl_tail:w #1#2 \q_stop

```

would give the wrong result for `\tl_tail:n { a { bc } }` (the braces would be stripped). Thus the only safe way to proceed is to first check that there is an item to grab (*i.e.* that the argument is not blank) and assuming there is to dispose of the first item. As with `\tl_head:n`, the result is protected from further expansion by `\unexpanded`. While we could optimise the test here, this would leave some tokens “banned” in the input, which we do not have with this definition.

```

4334   \exp_args:Nno \use:n { \cs_new:Npn \tl_tail:n #1 }
4335   {
4336     \exp_after:wN \_kernel_exp_not:w
4337     \tl_if_blank:nTF {#1}
4338     { { } }
4339     { \exp_after:wN { \use_none:n #1 } }
4340   }
4341   \cs_generate_variant:Nn \tl_tail:n { V , v , f }
4342   \cs_new:Npn \tl_tail:N { \exp_args:No \tl_tail:n }

```

(End definition for `\tl_head:N` and others. These functions are documented on page 54.)

`\tl_if_head_eq_meaning_p:nN` Accessing the first token of a token list is tricky in three cases: when it has category code 1 (begin-group token), when it is an explicit space, with category code 10 and character code 32, or when the token list is empty (obviously).

`\tl_if_head_eq_meaning:nN` Forgetting temporarily about this issue we would use the following test in `\tl_if_head_eq_charcode_p:nN`. Here, `\tl_head:w` yields the first token of the token list, then passed to `\exp_not:N`.

```

\if_charcode:w
  \exp_after:wN \exp_not:N \tl_head:w #1 \q_nil \q_stop
  \exp_not:N #2
  \_tl_head_exp_not:w
\_tl_if_head_eq_empty_arg:w

```

The two first special cases are detected by testing if the token list starts with an N-type token (the extra ? sends empty token lists to the true branch of this test). In those cases, the first token is a character, and since we only care about its character code, we can use `\str_head:n` to access it (this works even if it is a space character). An empty argument results in `\tl_head:w` leaving two tokens: `^` and `__tl_if_head_eq_empty_arg:w` which will result in the `\if_charcode:w` test being false and remove `\exp_not:N` and `#2`.

```

4343 \prg_new_conditional:Npnn \tl_if_head_eq_charcode:nN #1#2 { p , T , F , TF }
4344 {
4345   \if_charcode:w
4346     \tl_if_head_is_N_type:nTF { #1 ? }
4347     { \__tl_head_exp_not:w #1 { ^ \__tl_if_head_eq_empty_arg:w } \s__tl_stop }
4348     { \str_head:n {#1} }
4349     \exp_not:N #2
4350     \prg_return_true:
4351   \else:
4352     \prg_return_false:
4353   \fi:
4354 }
4355 \prg_generate_conditional_variant:Nnn \tl_if_head_eq_charcode:nN
4356 { f } { p , TF , T , F }

```

For `\tl_if_head_eq_catcode:nN`, again we detect special cases with a `\tl_if_head_is_N_type:n`. Then we need to test if the first token is a begin-group token or an explicit space token, and produce the relevant token, either `\c_group_begin_token` or `\c_space_token`. Again, for an empty argument, a hack is used, removing the token given by the user and leaving two tokens in the input stream which will make the `\if_catcode:w` test return false.

```

4357 \prg_new_conditional:Npnn \tl_if_head_eq_catcode:nN #1 #2 { p , T , F , TF }
4358 {
4359   \if_catcode:w
4360     \tl_if_head_is_N_type:nTF { #1 ? }
4361     { \__tl_head_exp_not:w #1 { ^ \__tl_if_head_eq_empty_arg:w } \s__tl_stop }
4362     {
4363       \tl_if_head_is_group:nTF {#1}
4364       \c_group_begin_token
4365       \c_space_token
4366     }
4367     \exp_not:N #2
4368     \prg_return_true:
4369   \else:
4370     \prg_return_false:
4371   \fi:
4372 }
4373 \prg_generate_conditional_variant:Nnn \tl_if_head_eq_catcode:nN
4374 { o } { p , TF , T , F }

```

For `\tl_if_head_eq_meaning:nN`, again, detect special cases. In the normal case, use `\tl_head:w`, with no `\exp_not:N` this time, since `\if_meaning:w` causes no expansion. With an empty argument, the test is true, and `\use_none:nnn` removes `#2` and `\prg_return_true:` and `\else:` (it is safe this way here as in this case `\prg_new_conditional:Npnn` didn't optimize these two away). In the special cases, we know that the first token is a character, hence `\if_charcode:w` and `\if_catcode:w` together are enough. We combine them in some order, hopefully faster than the reverse. Tests are

not nested because the arguments may contain unmatched primitive conditionals.

```

4375 \prg_new_conditional:Npnn \tl_if_head_eq_meaning:nN #1#2 { p , T , F , TF }
4376 {
4377   \tl_if_head_is_N_type:nTF { #1 ? }
4378     \__tl_if_head_eq_meaning_normal:nN
4379     \__tl_if_head_eq_meaning_special:nN
4380     {#1} #2
4381 }
4382 \cs_new:Npn \__tl_if_head_eq_meaning_normal:nN #1 #2
4383 {
4384   \exp_after:wN \if_meaning:w
4385     \__tl_tl_head:w #1 { ?? \use_none:nnn } \s__tl_stop #2
4386   \prg_return_true:
4387   \else:
4388     \prg_return_false:
4389   \fi:
4390 }
4391 \cs_new:Npn \__tl_if_head_eq_meaning_special:nN #1 #2
4392 {
4393   \if_charcode:w \str_head:n {#1} \exp_not:N #2
4394     \exp_after:wN \use_ii:nn
4395   \else:
4396     \prg_return_false:
4397   \fi:
4398   \use_none:n
4399   {
4400     \if_catcode:w \exp_not:N #2
4401       \tl_if_head_is_group:nTF {#1}
4402         { \c_group_begin_token }
4403         { \c_space_token }
4404     \prg_return_true:
4405   \else:
4406     \prg_return_false:
4407   \fi:
4408 }
4409 }

```

Both `\tl_if_head_eq_charcode:nN` and `\tl_if_head_eq_catcode:nN` will need to get the first token of their argument and apply `\exp_not:N` to it. `__tl_head_exp_not:w` does exactly that.

```

4410 \cs_new:Npn \__tl_head_exp_not:w #1 #2 \s__tl_stop
4411 { \exp_not:N #1 }

```

If the argument of `\tl_if_head_eq_charcode:nN` and `\tl_if_head_eq_catcode:nN` was empty `__tl_if_head_eq_empty_arg:w` will be left in the input stream. This macro has to remove `\exp_not:N` and the following token from the input stream to make sure no unbalanced if-construct is created and leave tokens there which make the two tests return false.

```

4412 \cs_new:Npn \__tl_if_head_eq_empty_arg:w \exp_not:N #1
4413 { ? }

```

(End definition for `\tl_if_head_eq_meaning:nNTF` and others. These functions are documented on page 55.)

`\tl_if_head_is_N_type_p:n`
`\tl_if_head_is_N_type:nTF`
`_tl_if_head_is_N_type_auxi:w`
`_tl_if_head_is_N_type_auxii:nn`
`_tl_if_head_is_N_type_auxiii:n`

A token list can be empty, can start with an explicit space character (catcode 10 and charcode 32), can start with a begin-group token (catcode 1), or start with an N-type argument. In the first two cases, the line involving `_tl_if_head_is_N_type_auxi:w` produces `f` (and otherwise nothing). In the third case (begin-group token), the lines involving `\token_to_str:N` produce a single closing brace. The category code test is thus true exactly in the fourth case, which is what we want. One cannot optimize by moving one of the `\scan_stop:` to the beginning: if `#1` contains primitive conditionals, all of its occurrences must be dealt with before the `\if:w` tries to skip the `true` branch of the conditional.

```

4414 \prg_new_conditional:Npnn \tl_if_head_is_N_type:n #1 { p , T , F , TF }
4415 {
4416   \if:w
4417     \if_false: { \fi: \_tl_if_head_is_N_type_auxi:w \prg_do_nothing: #1 ~ }
4418     { \exp_after:wN { \token_to_str:N #1 } }
4419     \scan_stop: \scan_stop:
4420     \prg_return_true:
4421   \else:
4422     \prg_return_false:
4423   \fi:
4424 }
4425 \exp_args:Nno \use:n { \cs_new:Npn \_tl_if_head_is_N_type_auxi:w #1 ~ }
4426 {
4427   \tl_if_empty:oTF { #1 }
4428   { f \exp_after:wN \use_none:nn }
4429   { \exp_after:wN \_tl_if_head_is_N_type_auxii:n }
4430   \exp_after:wN { \if_false: } \fi:
4431 }
4432 \cs_new:Npn \_tl_if_head_is_N_type_auxii:n #1
4433 { \exp_after:wN \use_none:n \exp_after:wN }

```

(End definition for `\tl_if_head_is_N_type:nTF` and others. This function is documented on page 55.)

`\tl_if_head_is_group_p:n`
`\tl_if_head_is_group:nTF`
`_tl_if_head_is_group_fi_false:w`

Pass the first token of `#1` through `\token_to_str:N`, then check for the brace balance. The extra `?` caters for an empty argument. This could be made faster, but we need all brace tricks to happen in one step of expansion, keeping the token list brace balanced at all times.

```

4434 \prg_new_conditional:Npnn \tl_if_head_is_group:n #1 { p , T , F , TF }
4435 {
4436   \if:w
4437     \exp_after:wN \use_none:n
4438     \exp_after:wN { \exp_after:wN { \token_to_str:N #1 ? } }
4439     \scan_stop: \scan_stop:
4440     \_tl_if_head_is_group_fi_false:w
4441   \fi:
4442   \if_true:
4443     \prg_return_true:
4444   \else:
4445     \prg_return_false:
4446   \fi:
4447 }
4448 \cs_new:Npn \_tl_if_head_is_group_fi_false:w \fi: \if_true: { \fi: \if_false: }

```

(End definition for `\tl_if_head_is_group:nTF` and `_tl_if_head_is_group_fi_false:w`. This function is documented on page 55.)

`\tl_if_head_is_space:p:n`
`\tl_if_head_is_space:nTF`
`__tl_if_head_is_space:w`

The auxiliary's argument is all that is before the first explicit space in `\prg_do_nothing:#1?~`. If that is a single `\prg_do_nothing:` the test yields `true`. Otherwise, that is more than one token, and the test yields `false`. The work is done within braces (with an `\if_false: { \fi: ... }` construction) both to hide potential alignment tab characters from \TeX in a table, and to allow for removing what remains of the token list after its first space. The use of `\if:w` ensures that the result of a single step of expansion directly yields a balanced token list (no trailing closing brace).

```

4449 \prg_new_conditional:Npnn \tl_if_head_is_space:n #1 { p , T , F , TF }
4450 {
4451   \if:w
4452     \if_false: { \fi: \__tl_if_head_is_space:w \prg_do_nothing: #1 ? ~ }
4453     \scan_stop: \scan_stop:
4454     \prg_return_true:
4455   \else:
4456     \prg_return_false:
4457   \fi:
4458 }
4459 \exp_args:Nno \use:n { \cs_new:Npn \__tl_if_head_is_space:w #1 ~ }
4460 {
4461   \__tl_if_empty_if:o {#1} \else: f \fi:
4462   \exp_after:wN \use_none:n \exp_after:wN { \if_false: } \fi:
4463 }

```

(End definition for `\tl_if_head_is_space:nTF` and `__tl_if_head_is_space:w`. This function is documented on page 55.)

8.12 Token by token changes

`\s__tl_act_stop`

The `__tl_act...` functions may be applied to any token list. Hence, we use a private quark, to allow any token, even quarks, in the token list. Only `\s__tl_act_stop` may not appear in the token lists manipulated by `__tl_act:NNNn` functions.

```

4464 \scan_new:N \s__tl_act_stop

```

(End definition for `\s__tl_act_stop`.)

`__tl_act:NNNn`
`__tl_act_output:n`
`__tl_act_reverse_output:n`
`__tl_act_loop:w`
`__tl_act_normal:NwNNN`
`__tl_act_group:nwNNN`
`__tl_act_space:wwNNN`
`__tl_act_end:w`
`\tl_act_if_head_is_space:nTF`
`__tl_act_if_head_is_space:w`
`\tl_act_if_head_is_space_true:w`
`\tl_use_none_delimit_by_q_act_stop:w`

To help control the expansion, `__tl_act:NNNn` should always be preceeded by `\exp:w` and ends by producing `\exp_end:` once the result has been obtained. This way no internal token of it can be accidentally end up in the input stream. Because `\s__tl_act_stop` can't appear without braces around it in the argument #1 of `__tl_act_loop:w`, we can use this marker to set up a fast test for leading spaces.

```

4465 \cs_set_protected:Npn \__tl_tmp:w #1
4466 {
4467   \cs_new:Npn \__tl_act_if_head_is_space:nTF ##1
4468   {
4469     \__tl_act_if_head_is_space:w
4470     \s__tl_act_stop ##1 \s__tl_act_stop \__tl_act_if_head_is_space_true:w
4471     \s__tl_act_stop #1 \s__tl_act_stop \use_ii:nn
4472   }
4473   \cs_new:Npn \__tl_act_if_head_is_space:w
4474   ##1 \s__tl_act_stop #1 ##2 \s__tl_act_stop
4475   {}
4476   \cs_new:Npn \__tl_act_if_head_is_space_true:w
4477   \s__tl_act_stop #1 \s__tl_act_stop \use_ii:nn ##1 ##2

```

```

4478     {##1}
4479   }
4480   \_tl\_tmp:w { ~ }

```

(We expand the definition `_tl_act_if_head_is_space:nTF` when setting up `_tl_act_loop:w`, so we can then undefine the auxiliary.) In the loop, we check how the token list begins and act accordingly. In the “group” case, we may have reached `\s_tl_act_stop`, the end of the list. Then leave `\exp_end:` and the result in the input stream, to terminate the expansion of `\exp:w`. Otherwise, apply the relevant function to the “arguments”, #3 and to the head of the token list. Then repeat the loop. The scheme is the same if the token list starts with an N-type or with a space, making sure that `_tl_act_space:wwNNN` gobbles the space.

```

4481 \exp\_args:Nnx \use:n { \cs\_new:Npn \_tl\_act\_loop:w #1 \s\_tl\_act\_stop }
4482 {
4483   \exp\_not:o { \_tl\_act\_if\_head\_is\_space:nTF {#1} }
4484   \exp\_not:N \_tl\_act\_space:wwNNN
4485   {
4486     \exp\_not:o { \tl\_if\_head\_is\_group:nTF {#1} }
4487     \exp\_not:N \_tl\_act\_group:nwNNN
4488     \exp\_not:N \_tl\_act\_normal:NwNNN
4489   }
4490   \exp\_not:n {#1} \s\_tl\_act\_stop
4491 }
4492 \cs\_undefine:N \_tl\_act\_if\_head\_is\_space:nTF
4493 \cs\_new:Npn \_tl\_act\_normal:NwNNN #1 #2 \s\_tl\_act\_stop #3
4494 {
4495   #3 #1
4496   \_tl\_act\_loop:w #2 \s\_tl\_act\_stop
4497   #3
4498 }
4499 \cs\_new:Npn \_tl\_use\_none\_delimit\_by\_s\_act\_stop:w #1 \s\_tl\_act\_stop { }
4500 \cs\_new:Npn \_tl\_act\_end:wn #1 \_tl\_act\_result:n #2
4501 { \group\_align\_safe\_end: \exp\_end: #2 }
4502 \cs\_new:Npn \_tl\_act\_group:nwNNN #1 #2 \s\_tl\_act\_stop #3#4#5
4503 {
4504   \_tl\_use\_none\_delimit\_by\_s\_act\_stop:w #1 \_tl\_act\_end:wn \s\_tl\_act\_stop
4505   #5 {#1}
4506   \_tl\_act\_loop:w #2 \s\_tl\_act\_stop
4507   #3 #4 #5
4508 }
4509 \exp\_last\_unbraced:NNo
4510 \cs\_new:Npn \_tl\_act\_space:wwNNN \c\_space\_tl #1 \s\_tl\_act\_stop #2#3
4511 {
4512   #3
4513   \_tl\_act\_loop:w #1 \s\_tl\_act\_stop
4514   #2 #3
4515 }

```

`_tl_act:NNNn` loops over tokens, groups, and spaces in #4. `{\s_@@_act_stop}` serves as the end of token list marker, the ? after it avoids losing outer braces. The result is stored as an argument for the dummy function `_tl_act_result:n`.

```

4516 \cs\_new:Npn \_tl\_act:NNNn #1#2#3#4
4517 {
4518   \group\_align\_safe\_begin:

```

```

4519     \_tl_act_loop:w #4 { \s\_tl_act_stop } ? \s\_tl_act_stop
4520     #1 #3 #2
4521     \_tl_act_result:n { }
4522 }

```

Typically, the output is done to the right of what was already output, using `_tl_act_output:n`, but for the `_tl_act_reverse` functions, it should be done to the left.

```

4523 \cs_new:Npn \_tl_act_output:n #1 #2 \_tl_act_result:n #3
4524 { #2 \_tl_act_result:n { #3 #1 } }
4525 \cs_new:Npn \_tl_act_reverse_output:n #1 #2 \_tl_act_result:n #3
4526 { #2 \_tl_act_result:n { #1 #3 } }

```

(End definition for `_tl_act:NNNn` and others.)

`\tl_reverse:n`
`\tl_reverse:o`
`\tl_reverse:V`

The goal here is to reverse without losing spaces nor braces. This is done using the general internal function `_tl_act:NNNn`. Spaces and “normal” tokens are output on the left of the current output. Grouped tokens are output to the left but without any reversal within the group.

```

\_tl_reverse_normal:nN
\_tl_reverse_group_preserve:nN
\_tl_reverse_space:n
4527 \cs_new:Npn \tl_reverse:n #1
4528 {
4529     \_kernel_exp_not:w \exp_after:wN
4530     {
4531         \exp:w
4532         \_tl_act:NNNn
4533         \_tl_reverse_normal:N
4534         \_tl_reverse_group_preserve:n
4535         \_tl_reverse_space:
4536         {#1}
4537     }
4538 }
4539 \cs_generate_variant:Nn \tl_reverse:n { o , V }
4540 \cs_new:Npn \_tl_reverse_normal:N
4541 { \_tl_act_reverse_output:n }
4542 \cs_new:Npn \_tl_reverse_group_preserve:n #1
4543 { \_tl_act_reverse_output:n { {#1} } }
4544 \cs_new:Npn \_tl_reverse_space:
4545 { \_tl_act_reverse_output:n { ~ } }

```

(End definition for `\tl_reverse:n` and others. This function is documented on page 52.)

`\tl_reverse:N`
`\tl_reverse:c`
`\tl_greverse:N`
`\tl_greverse:c`

This reverses the list, leaving `\exp_stop_f:` in front, which stops the `f`-expansion.

```

4546 \cs_new_protected:Npn \tl_reverse:N #1
4547 { \_kernel_tl_set:Nx #1 { \exp_args:No \tl_reverse:n { #1 } } }
4548 \cs_new_protected:Npn \tl_greverse:N #1
4549 { \_kernel_tl_gset:Nx #1 { \exp_args:No \tl_reverse:n { #1 } } }
4550 \cs_generate_variant:Nn \tl_reverse:N { c }
4551 \cs_generate_variant:Nn \tl_greverse:N { c }

```

(End definition for `\tl_reverse:N` and `\tl_greverse:N`. These functions are documented on page 52.)

8.13 Using a single item

`\tl_item:nn` The idea here is to find the offset of the item from the left, then use a loop to grab the correct item. If the resulting offset is too large, then `__tl_if_recursion_tail_break:nN` terminates the loop, and returns nothing at all.

`\tl_item:Nn`

`\tl_item:cn`

`__tl_item_aux:nn`

`__tl_item:nn`

```

4552 \cs_new:Npn \tl_item:nn #1#2
4553 {
4554   \exp_args:Nf \__tl_item:nn
4555   { \exp_args:Nf \__tl_item_aux:nn { \int_eval:n {#2} } {#1} }
4556   #1
4557   \q_tl_recursion_tail
4558   \prg_break_point:
4559 }
4560 \cs_new:Npn \__tl_item_aux:nn #1#2
4561 {
4562   \int_compare:nNnTF {#1} < 0
4563   { \int_eval:n { \tl_count:n {#2} + 1 + #1 } }
4564   {#1}
4565 }
4566 \cs_new:Npn \__tl_item:nn #1#2
4567 {
4568   \__tl_if_recursion_tail_break:nN {#2} \prg_break:
4569   \int_compare:nNnTF {#1} = 1
4570   { \prg_break:n { \exp_not:n {#2} } }
4571   { \exp_args:Nf \__tl_item:nn { \int_eval:n { #1 - 1 } } }
4572 }
4573 \cs_new:Npn \tl_item:Nn { \exp_args:No \tl_item:nn }
4574 \cs_generate_variant:Nn \tl_item:Nn { c }

```

(End definition for `\tl_item:nn` and others. These functions are documented on page 56.)

`\tl_rand_item:n` Importantly `\tl_item:nn` only evaluates its argument once.

`\tl_rand_item:N`

`\tl_rand_item:c`

```

4575 \cs_new:Npn \tl_rand_item:n #1
4576 {
4577   \tl_if_blank:nF {#1}
4578   { \tl_item:nn {#1} { \int_rand:nn { 1 } { \tl_count:n {#1} } } }
4579 }
4580 \cs_new:Npn \tl_rand_item:N { \exp_args:No \tl_rand_item:n }
4581 \cs_generate_variant:Nn \tl_rand_item:N { c }

```

(End definition for `\tl_rand_item:n` and `\tl_rand_item:N`. These functions are documented on page 56.)

`\tl_range:Nnn` To avoid checking for the end of the token list at every step, start by counting the number l of items and “normalizing” the bounds, namely clamping them to the interval $[0, l]$ and dealing with negative indices. More precisely, `__tl_range_items:nnNn` receives the number of items to skip at the beginning of the token list, the index of the last item to keep, a function which is either `__tl_range:w` or the token list itself. If nothing should be kept, leave `{}`: this stops the f-expansion of `\tl_head:f` and that function produces an empty result. Otherwise, repeatedly call `__tl_range_skip:w` to delete $\#1$ items from the input stream (the extra brace group avoids an off-by-one shift). For the braced version `__tl_range_braced:w` sets up `__tl_range_collect_braced:w` which stores items one by one in an argument after the semicolon. Depending on the first token of the tail, either just move it (if it is a space) or also decrement the number of items left

`__tl_range:Nnnn`

`__tl_range:nnnNn`

`__tl_range:nnNn`

`__tl_range_skip:w`

`__tl_range:w`

`__tl_range_skip_spaces:n`

`__tl_range_collect:nn`

`__tl_range_collect:ff`

`__tl_range_collect_space:nw`

`__tl_range_collect_N:nN`

`__tl_range_collect_group:nN`

to find. Eventually, the result is a brace group followed by the rest of the token list, and `\tl_head:f` cleans up and gives the result in `\exp_not:n`.

```

4582 \cs_new:Npn \tl_range:Nnn { \exp_args:No \tl_range:nnn }
4583 \cs_generate_variant:Nn \tl_range:Nnn { c }
4584 \cs_new:Npn \tl_range:nnn { \__tl_range:Nnnn \__tl_range:w }
4585 \cs_new:Npn \__tl_range:Nnnn #1#2#3#4
4586 {
4587   \tl_head:f
4588   {
4589     \exp_args:Nf \__tl_range:nnnNn
4590     { \tl_count:n {#2} } {#3} {#4} #1 {#2}
4591   }
4592 }
4593 \cs_new:Npn \__tl_range:nnnNn #1#2#3
4594 {
4595   \exp_args:Nff \__tl_range:nnNn
4596   {
4597     \exp_args:Nf \__tl_range_normalize:nn
4598     { \int_eval:n { #2 - 1 } } {#1}
4599   }
4600   {
4601     \exp_args:Nf \__tl_range_normalize:nn
4602     { \int_eval:n {#3} } {#1}
4603   }
4604 }
4605 \cs_new:Npn \__tl_range:nnNn #1#2#3#4
4606 {
4607   \if_int_compare:w #2 > #1 \exp_stop_f: \else:
4608     \exp_after:wN { \exp_after:wN }
4609   \fi:
4610   \exp_after:wN #3
4611   \int_value:w \int_eval:n { #2 - #1 } \exp_after:wN ;
4612   \exp_after:wN { \exp:w \__tl_range_skip:w #1 ; { } #4 }
4613 }
4614 \cs_new:Npn \__tl_range_skip:w #1 ; #2
4615 {
4616   \if_int_compare:w #1 > 0 \exp_stop_f:
4617     \exp_after:wN \__tl_range_skip:w
4618     \int_value:w \int_eval:n { #1 - 1 } \exp_after:wN ;
4619   \else:
4620     \exp_after:wN \exp_end:
4621   \fi:
4622 }
4623 \cs_new:Npn \__tl_range:w #1 ; #2
4624 {
4625   \exp_args:Nf \__tl_range_collect:nn
4626   { \__tl_range_skip_spaces:n {#2} } {#1}
4627 }
4628 \cs_new:Npn \__tl_range_skip_spaces:n #1
4629 {
4630   \tl_if_head_is_space:nTF {#1}
4631   { \exp_args:Nf \__tl_range_skip_spaces:n {#1} }
4632   { { } #1 }
4633 }

```

```

4634 \cs_new:Npn \__tl_range_collect:nn #1#2
4635 {
4636   \int_compare:nNnTF {#2} = 0
4637   {#1}
4638   {
4639     \exp_args:No \tl_if_head_is_space:nTF { \use_none:n #1 }
4640     {
4641       \exp_args:Nf \__tl_range_collect:nn
4642       { \__tl_range_collect_space:nw #1 }
4643       {#2}
4644     }
4645     {
4646       \__tl_range_collect:ff
4647       {
4648         \exp_args:No \tl_if_head_is_N_type:nTF { \use_none:n #1 }
4649         { \__tl_range_collect_N:nN }
4650         { \__tl_range_collect_group:nn }
4651         #1
4652       }
4653       { \int_eval:n { #2 - 1 } }
4654     }
4655   }
4656 }
4657 \cs_new:Npn \__tl_range_collect_space:nw #1 ~ { { #1 ~ } }
4658 \cs_new:Npn \__tl_range_collect_N:nN #1#2 { { #1 #2 } }
4659 \cs_new:Npn \__tl_range_collect_group:nn #1#2 { { #1 {#2} } }
4660 \cs_generate_variant:Nn \__tl_range_collect:nn { ff }

```

(End definition for `\tl_range:Nnn` and others. These functions are documented on page 57.)

`__tl_range_normalize:nn` This function converts an $\langle index \rangle$ argument into an explicit position in the token list (a result of 0 denoting “out of bounds”). Expects two explicit integer arguments: the $\langle index \rangle$ #1 and the string count #2. If #1 is negative, replace it by #1 + #2 + 1, then limit to the range [0, #2].

```

4661 \cs_new:Npn \__tl_range_normalize:nn #1#2
4662 {
4663   \int_eval:n
4664   {
4665     \if_int_compare:w #1 < 0 \exp_stop_f:
4666     \if_int_compare:w #1 < -#2 \exp_stop_f:
4667     0
4668     \else:
4669     #1 + #2 + 1
4670     \fi:
4671     \else:
4672     \if_int_compare:w #1 < #2 \exp_stop_f:
4673     #1
4674     \else:
4675     #2
4676     \fi:
4677     \fi:
4678   }
4679 }

```

(End definition for `__tl_range_normalize:nn`.)

8.14 Viewing token lists

`\tl_show:N` Showing token list variables is done after checking that the variable is defined (see `__tl_show:c` kernel_register_show:N).

`\tl_log:N` 4680 \cs_new_protected:Npn \tl_show:N { __tl_show:NN \tl_show:n }

`\tl_log:c` 4681 \cs_generate_variant:Nn \tl_show:N { c }

`__tl_show:NN` 4682 \cs_new_protected:Npn \tl_log:N { __tl_show:NN \tl_log:n }

4683 \cs_generate_variant:Nn \tl_log:N { c }

4684 \cs_new_protected:Npn __tl_show:NN #1#2

4685 {

4686 __kernel_chk_defined:NT #2

4687 {

4688 \exp_args:Ne #1

4689 { \token_to_str:N #2 = __kernel_exp_not:w \exp_after:wN {#2} }

4690 }

4691 }

(End definition for `\tl_show:N`, `\tl_log:N`, and `__tl_show:NN`. These functions are documented on page 58.)

`\tl_show:n` Many show functions are based on `\tl_show:n`. The argument of `\tl_show:n` is line-wrapped using `\iow_wrap:nnnN` but with a leading `>~` and trailing period, both removed before passing the wrapped text to the `\showtokens` primitive. This primitive shows the result with a leading `>~` and trailing period.

`__tl_show:n`

`__tl_show:w`

The token list `\l__tl_internal_a_tl` containing the result of all these manipulations is displayed to the terminal using `\tex_showtokens:D` and an odd `\exp_after:wN` which expand the closing brace to improve the output slightly. The calls to `__kernel_iow_with:Nnn` ensure that the `\newlinechar` is set to 10 so that the `\iow_newline:` inserted by the line-wrapping code are correctly recognized by T_EX, and that `\errorcontextlines` is `-1` to avoid printing irrelevant context.

4692 \cs_new_protected:Npn \tl_show:n #1

4693 { \iow_wrap:nnnN { >~ \tl_to_str:n {#1} . } { } { } __tl_show:n }

4694 \cs_new_protected:Npn __tl_show:n #1

4695 {

4696 \tl_set:Nf \l__tl_internal_a_tl { __tl_show:w #1 \s__tl_stop }

4697 __kernel_iow_with:Nnn \tex_newlinechar:D { 10 }

4698 {

4699 __kernel_iow_with:Nnn \tex_errorcontextlines:D { -1 }

4700 {

4701 \tex_showtokens:D \exp_after:wN \exp_after:wN \exp_after:wN

4702 { \exp_after:wN \l__tl_internal_a_tl }

4703 }

4704 }

4705 }

4706 \cs_new:Npn __tl_show:w #1 > #2 . \s__tl_stop {#2}

(End definition for `\tl_show:n`, `__tl_show:n`, and `__tl_show:w`. This function is documented on page 58.)

`\tl_log:n` Logging is much easier, simply line-wrap. The `>~` and trailing period is there to match the output of `\tl_show:n`.

4707 \cs_new_protected:Npn \tl_log:n #1

4708 { \iow_wrap:nnnN { > ~ \tl_to_str:n {#1} . } { } { } \iow_log:n }

(End definition for `\tl_log:n`. This function is documented on page 58.)

8.15 Internal scan marks

`\s__tl_nil` Internal scan marks. These are defined here at the end because the code for `\scan_new:N` depends on some `l3tl` functions.

`\s__tl_mark`

`\s__tl_stop`

```

4709 \scan_new:N \s__tl_nil
4710 \scan_new:N \s__tl_mark
4711 \scan_new:N \s__tl_stop

```

(End definition for `\s__tl_nil`, `\s__tl_mark`, and `\s__tl_stop`.)

8.16 Scratch token lists

`\g_tmpa_tl` Global temporary token list variables. They are supposed to be set and used immediately, with no delay between the definition and the use because you can't count on other macros not to redefine them from under you.

`\g_tmpb_tl`

```

4712 \tl_new:N \g_tmpa_tl
4713 \tl_new:N \g_tmpb_tl

```

(End definition for `\g_tmpa_tl` and `\g_tmpb_tl`. These variables are documented on page 59.)

`\l_tmpa_tl` These are local temporary token list variables. Be sure not to assume that the value you put into them will survive for long—see discussion above.

`\l_tmpb_tl`

```

4714 \tl_new:N \l_tmpa_tl
4715 \tl_new:N \l_tmpb_tl

```

(End definition for `\l_tmpa_tl` and `\l_tmpb_tl`. These variables are documented on page 59.)

We finally clean up a temporary control sequence that we have used at various points to set up some definitions.

```

4716 \cs_undefine:N \__tl_tmp:w
4717 \</package>

```

9 l3str implementation

```

4718 \*package>
4719 \@@=str>

```

9.1 Internal auxiliaries

`\s__str_mark` Internal scan marks.

`\s__str_stop`

```

4720 \scan_new:N \s__str_mark
4721 \scan_new:N \s__str_stop

```

(End definition for `\s__str_mark` and `\s__str_stop`.)

`__str_use_none_delimit_by_s_stop:w` Functions to gobble up to a scan mark.

`__str_use_i_delimit_by_s_stop:nw`

```

4722 \cs_new:Npn \__str_use_none_delimit_by_s_stop:w #1 \s__str_stop { }
4723 \cs_new:Npn \__str_use_i_delimit_by_s_stop:nw #1 #2 \s__str_stop {#1}

```

(End definition for `__str_use_none_delimit_by_s_stop:w` and `__str_use_i_delimit_by_s_stop:nw`.)

`\q__str_recursion_tail` Internal recursion quarks.

`\q__str_recursion_stop`

```

4724 \quark_new:N \q__str_recursion_tail
4725 \quark_new:N \q__str_recursion_stop

```

(End definition for `\q__str_recursion_tail` and `\q__str_recursion_stop`.)

`__str_if_recursion_tail_break:NN`
`__str_if_recursion_tail_stop_do:Nn`

Functions to query recursion quarks.

```
4726 \__kernel_quark_new_test:N \__str_if_recursion_tail_break:NN
4727 \__kernel_quark_new_test:N \__str_if_recursion_tail_stop_do:Nn
```

(End definition for `__str_if_recursion_tail_break:NN` and `__str_if_recursion_tail_stop_do:Nn`.)

9.2 Creating and setting string variables

`\str_new:N` A string is simply a token list. The full mapping system isn't set up yet so do things by hand.
`\str_new:c`
`\str_use:N` 4728 `\group_begin:`
`\str_use:c` 4729 `\cs_set_protected:Npn __str_tmp:n #1`
`\str_clear:N` 4730 `{`
`\str_clear:c` 4731 `\tl_if_blank:nF {#1}`
`\str_gclear:N` 4732 `{`
`\str_gclear:c` 4733 `\cs_new_eq:cc { str_ #1 :N } { tl_ #1 :N }`
`\str_clear_new:N` 4734 `\exp_args:Nc \cs_generate_variant:Nn { str_ #1 :N } { c }`
`\str_clear_new:c` 4735 `__str_tmp:n`
`\str_gclear_new:N` 4736 `}`
`\str_gclear_new:c` 4737 `}`
`\str_set_eq:NN` 4738 `__str_tmp:n`
`\str_set_eq:cN` 4739 `{ new }`
`\str_set_eq:Nc` 4740 `{ use }`
`\str_set_eq:cc` 4741 `{ clear }`
`\str_gset_eq:NN` 4742 `{ gclear }`
`\str_gset_eq:cN` 4743 `{ clear_new }`
`\str_gset_eq:Nc` 4744 `{ gclear_new }`
`\str_gset_eq:cc` 4745 `{ }`
`\str_concat:NNN` 4746 `\group_end:`
`\str_concat:ccc` 4747 `\cs_new_eq:NN \str_set_eq:NN \tl_set_eq:NN`
`\str_gconcat:NNN` 4748 `\cs_new_eq:NN \str_gset_eq:NN \tl_gset_eq:NN`
`\str_gconcat:ccc` 4749 `\cs_generate_variant:Nn \str_set_eq:NN { c , Nc , cc }`
4750 `\cs_generate_variant:Nn \str_gset_eq:NN { c , Nc , cc }`
4751 `\cs_new_eq:NN \str_concat:NNN \tl_concat:NNN`
4752 `\cs_new_eq:NN \str_gconcat:NNN \tl_gconcat:NNN`
4753 `\cs_generate_variant:Nn \str_concat:NNN { ccc }`
4754 `\cs_generate_variant:Nn \str_gconcat:NNN { ccc }`

(End definition for `\str_new:N` and others. These functions are documented on page 60.)

`\str_set:Nn` Simply convert the token list inputs to $\langle strings \rangle$.
`\str_set:NV` 4755 `\group_begin:`
`\str_set:Nx` 4756 `\cs_set_protected:Npn __str_tmp:n #1`
`\str_set:cn` 4757 `{`
`\str_set:cV` 4758 `\tl_if_blank:nF {#1}`
`\str_set:cx` 4759 `{`
`\str_gset:Nn` 4760 `\cs_new_protected:cpx { str_ #1 :Nn } ##1##2`
`\str_gset:NV` 4761 `{`
`\str_gset:Nx` 4762 `\exp_not:c { tl_ #1 :Nx } ##1`
`\str_gset:cn` 4763 `{ \exp_not:N \tl_to_str:n {##2} }`
`\str_gset:cV` 4764 `}`
`\str_gset:cx` 4765 `\cs_generate_variant:cn { str_ #1 :Nn } { NV , Nx , cn , cV , cx }`
`\str_const:Nn`
`\str_const:NV`
`\str_const:Nx`
`\str_const:cn`
`\str_const:cV`
`\str_const:cx`
`\str_put_left:Nn`
`\str_put_left:NV`
`\str_put_left:Nx`

```

4766         \__str_tmp:n
4767     }
4768 }
4769 \__str_tmp:n
4770 { set }
4771 { gset }
4772 { const }
4773 { put_left }
4774 { gput_left }
4775 { put_right }
4776 { gput_right }
4777 { }
4778 \group_end:

```

(End definition for `\str_set:Nn` and others. These functions are documented on page 61.)

9.3 Modifying string variables

`\str_replace_all:Nnn` Start by applying `\tl_to_str:n` to convert the old and new token lists to strings, and also apply `\tl_to_str:N` to avoid any issues if we are fed a token list variable. Then
`\str_replace_all:cnn` the code is a much simplified version of the token list code because neither the delimiter
`\str_greplace_all:Nnn` nor the replacement can contain macro parameters or braces. The delimiter `\s__str_`
`\str_greplace_all:cnn` mark cannot appear in the string to edit so it is used in all cases. Some x-expansion is
`\str_replace_once:Nnn` unnecessary. There is no need to avoid losing braces nor to protect against expansion.
`\str_replace_once:cnn` The ending code is much simplified and does not need to hide in braces.
`\str_greplace_once:Nnn`
`\str_greplace_once:cnn`

```

4779 \cs_new_protected:Npn \str_replace_once:Nnn
4780 { \__str_replace:NNNnn \prg_do_nothing: \__kernel_tl_set:Nx }
4781 \cs_new_protected:Npn \str_greplace_once:Nnn
4782 { \__str_replace:NNNnn \prg_do_nothing: \__kernel_tl_gset:Nx }
4783 \cs_new_protected:Npn \str_replace_all:Nnn
4784 { \__str_replace:NNNnn \__str_replace_next:w \__kernel_tl_set:Nx }
4785 \cs_new_protected:Npn \str_greplace_all:Nnn
4786 { \__str_replace:NNNnn \__str_replace_next:w \__kernel_tl_gset:Nx }
4787 \cs_generate_variant:Nn \str_replace_once:Nnn { c }
4788 \cs_generate_variant:Nn \str_greplace_once:Nnn { c }
4789 \cs_generate_variant:Nn \str_replace_all:Nnn { c }
4790 \cs_generate_variant:Nn \str_greplace_all:Nnn { c }
4791 \cs_new_protected:Npn \__str_replace:NNNnn #1#2#3#4#5
4792 {
4793     \tl_if_empty:nTF {#4}
4794     {
4795         \__kernel_msg_error:nxx { kernel } { empty-search-pattern } {#5}
4796     }
4797     {
4798         \use:x
4799         {
4800             \exp_not:n { \__str_replace_aux:NNNnnn #1 #2 #3 }
4801             { \tl_to_str:N #3 }
4802             { \tl_to_str:n {#4} } { \tl_to_str:n {#5} }
4803         }
4804     }
4805 }
4806 \cs_new_protected:Npn \__str_replace_aux:NNNnnn #1#2#3#4#5#6

```

```

4807 {
4808   \cs_set:Npn \__str_replace_next:w ##1 #5 { ##1 #6 #1 }
4809   #2 #3
4810   {
4811     \__str_replace_next:w
4812     #4
4813     \__str_use_none_delimit_by_s_stop:w
4814     #5
4815     \s__str_stop
4816   }
4817 }
4818 \cs_new_eq:NN \__str_replace_next:w ?

```

(End definition for `\str_replace_all:Nnn` and others. These functions are documented on page 62.)

```

\str_remove_once:Nn Removal is just a special case of replacement.
\str_remove_once:cn 4819 \cs_new_protected:Npn \str_remove_once:Nn #1#2
\str_gremove_once:Nn 4820 { \str_replace_once:Nnn #1 {#2} { } }
\str_gremove_once:cn 4821 \cs_new_protected:Npn \str_gremove_once:Nn #1#2
4822 { \str_greplace_once:Nnn #1 {#2} { } }
4823 \cs_generate_variant:Nn \str_remove_once:Nn { c }
4824 \cs_generate_variant:Nn \str_gremove_once:Nn { c }

```

(End definition for `\str_remove_once:Nn` and `\str_gremove_once:Nn`. These functions are documented on page 62.)

```

\str_remove_all:Nn Removal is just a special case of replacement.
\str_remove_all:cn 4825 \cs_new_protected:Npn \str_remove_all:Nn #1#2
\str_gremove_all:Nn 4826 { \str_replace_all:Nnn #1 {#2} { } }
\str_gremove_all:cn 4827 \cs_new_protected:Npn \str_gremove_all:Nn #1#2
4828 { \str_greplace_all:Nnn #1 {#2} { } }
4829 \cs_generate_variant:Nn \str_remove_all:Nn { c }
4830 \cs_generate_variant:Nn \str_gremove_all:Nn { c }

```

(End definition for `\str_remove_all:Nn` and `\str_gremove_all:Nn`. These functions are documented on page 62.)

9.4 String comparisons

```

\str_if_empty_p:N More copy-paste!
\str_if_empty_p:c 4831 \prg_new_eq_conditional:NNn \str_if_exist:N \tl_if_exist:N
\str_if_empty:NTF 4832 { p , T , F , TF }
\str_if_empty:cTF 4833 \prg_new_eq_conditional:NNn \str_if_exist:c \tl_if_exist:c
\str_if_exist_p:N 4834 { p , T , F , TF }
\str_if_exist_p:c 4835 \prg_new_eq_conditional:NNn \str_if_empty:N \tl_if_empty:N
\str_if_exist:NTF 4836 { p , T , F , TF }
\str_if_exist:cTF 4837 \prg_new_eq_conditional:NNn \str_if_empty:c \tl_if_empty:c
4838 { p , T , F , TF }

```

(End definition for `\str_if_empty:NTF` and `\str_if_exist:NTF`. These functions are documented on page 63.)

```

\__str_if_eq:nn String comparisons rely on the primitive \(pdf)strcmp, so we define a new name for it.
4839 \cs_new_eq:NN \__str_if_eq:nn \tex_strcmp:D

```

(End definition for `_str_if_eq:nn`.)

`\str_if_eq_p:nn` Modern engines provide a direct way of comparing two token lists, but returning a number. This set of conditionals therefore make life a bit clearer. The `nn` and `xx` versions are created directly as this is most efficient.

```

4840 \prg_new_conditional:Npnn \str_if_eq:nn #1#2 { p , T , F , TF }
4841 {
4842   \if_int_compare:w
4843     \_str_if_eq:nn { \exp_not:n {#1} } { \exp_not:n {#2} }
4844     = 0 \exp_stop_f:
4845     \prg_return_true: \else: \prg_return_false: \fi:
4846 }
4847 \prg_generate_conditional_variant:Nnn \str_if_eq:nn
4848 { V , v , o , nV , no , VV , nv } { p , T , F , TF }
4849 \prg_new_conditional:Npnn \str_if_eq:ee #1#2 { p , T , F , TF }
4850 {
4851   \if_int_compare:w \_str_if_eq:nn {#1} {#2} = 0 \exp_stop_f:
4852   \prg_return_true: \else: \prg_return_false: \fi:
4853 }

```

(End definition for `\str_if_eq:nnTF`. This function is documented on page 63.)

`\str_if_eq_p:NN` Note that `\str_if_eq:NN` is different from `\tl_if_eq:NN` because it needs to ignore category codes.

```

4854 \prg_new_conditional:Npnn \str_if_eq:NN #1#2 { p , TF , T , F }
4855 {
4856   \if_int_compare:w
4857     \_str_if_eq:nn { \tl_to_str:N #1 } { \tl_to_str:N #2 }
4858     = 0 \exp_stop_f: \prg_return_true: \else: \prg_return_false: \fi:
4859 }
4860 \prg_generate_conditional_variant:Nnn \str_if_eq:NN
4861 { c , Nc , cc } { T , F , TF , p }

```

(End definition for `\str_if_eq:NNTF`. This function is documented on page 63.)

`\str_if_in:NnTF` Everything here needs to be detokenized but beyond that it is a simple token list test.
`\str_if_in:cnTF` It would be faster to fine-tune the T, F, TF variants by calling the appropriate variant of
`\str_if_in:nnTF` `\tl_if_in:nnTF` directly but that takes more code.

```

4862 \prg_new_protected_conditional:Npnn \str_if_in:Nn #1#2 { T , F , TF }
4863 {
4864   \use:x
4865   { \tl_if_in:nnTF { \tl_to_str:N #1 } { \tl_to_str:n {#2} } }
4866   { \prg_return_true: } { \prg_return_false: }
4867 }
4868 \prg_generate_conditional_variant:Nnn \str_if_in:Nn
4869 { c } { T , F , TF }
4870 \prg_new_protected_conditional:Npnn \str_if_in:nn #1#2 { T , F , TF }
4871 {
4872   \use:x
4873   { \tl_if_in:nnTF { \tl_to_str:n {#1} } { \tl_to_str:n {#2} } }
4874   { \prg_return_true: } { \prg_return_false: }
4875 }

```

(End definition for `\str_if_in:NnTF` and `\str_if_in:nnTF`. These functions are documented on page 63.)

`\str_case:nn` Much the same as `\tl_case:nn(TF)` here: just a change in the internal comparison.

```

\str_case:Vn 4876 \cs_new:Npn \str_case:nn #1#2
\str_case:on 4877 {
\str_case:nV 4878   \exp:w
\str_case:nv 4879   \__str_case:nnTF {#1} {#2} { } { }
4880 }
\str_case:nnTF 4881 \cs_new:Npn \str_case:nnT #1#2#3
\str_case:VnTF 4882 {
\str_case:onTF 4883   \exp:w
\str_case:nVTF 4884   \__str_case:nnTF {#1} {#2} {#3} { }
\str_case:nvTF 4885   }
\str_case_e:nn 4886 \cs_new:Npn \str_case:nnF #1#2
\str_case_e:nnTF 4887 {
  \__str_case:nnTF 4888   \exp:w
  \__str_case_e:nnTF 4889   \__str_case:nnTF {#1} {#2} { }
  \__str_case:nw 4890   }
  \__str_case_e:nw 4891 \cs_new:Npn \str_case:nnTF #1#2
  \__str_case_end:nw 4892 {
    \exp:w 4893
    \__str_case:nnTF {#1} {#2} 4894
  } 4895
  \cs_new:Npn \__str_case:nnTF #1#2#3#4 4896
  { \__str_case:nw {#1} #2 {#1} { } \s__str_mark {#3} \s__str_mark {#4} \s__str_stop } 4897
  \cs_generate_variant:Nn \str_case:nn { V , o , nV , nv } 4898
  \prg_generate_conditional_variant:Nnn \str_case:nn 4899
  { V , o , nV , nv } { T , F , TF } 4900
  \cs_new:Npn \__str_case:nw #1#2#3 4901
  {
    \str_if_eq:nnTF {#1} {#2} 4902
    { \__str_case_end:nw {#3} } 4903
    { \__str_case:nw {#1} } 4904
  } 4905
  \cs_new:Npn \str_case_e:nn #1#2 4906
  {
    \exp:w 4907
    \__str_case_e:nnTF {#1} {#2} { } { } 4908
  } 4909
  \cs_new:Npn \str_case_e:nnT #1#2#3 4910
  {
    \exp:w 4911
    \__str_case_e:nnTF {#1} {#2} {#3} { } 4912
  } 4913
  \cs_new:Npn \str_case_e:nnF #1#2 4914
  {
    \exp:w 4915
    \__str_case_e:nnTF {#1} {#2} { } 4916
  } 4917
  \cs_new:Npn \str_case_e:nnTF #1#2 4918
  {
    \exp:w 4919
    \__str_case_e:nnTF {#1} {#2} 4920
  } 4921
  \cs_new:Npn \__str_case_e:nnTF #1#2#3#4 4922
  { \__str_case_e:nw {#1} #2 {#1} { } \s__str_mark {#3} \s__str_mark {#4} \s__str_stop } 4923

```

```

4929 \cs_new:Npn \__str_case_e:nw #1#2#3
4930 {
4931   \str_if_eq:eeTF {#1} {#2}
4932     { \__str_case_end:nw {#3} }
4933     { \__str_case_e:nw {#1} }
4934 }
4935 \cs_new:Npn \__str_case_end:nw #1#2#3 \s__str_mark #4#5 \s__str_stop
4936 { \exp_end: #1 #4 }

```

(End definition for `\str_case:nnTF` and others. These functions are documented on page 64.)

9.5 Mapping to strings

`\str_map_function:NN` The inline and variable mappings are similar to the usual token list mappings but start out by turning the argument to an “other string”. Doing the same for the expandable function mapping would require `__kernel_str_to_other:n`, quadratic in the string length. To deal with spaces in that case, `__str_map_function:w` replaces the following space by a braced space and a further call to itself. These are received by `__str_map_function:Nn`, which passes the space to #1 and calls `__str_map_function:w` to deal with the next space. The space before the braced space allows to optimize the `\q__str_recursion_tail` test. Of course we need to include a trailing space (the question mark is needed to avoid losing the space when TeX tokenizes the line). At the cost of about three more auxiliaries this code could get a 9 times speed up by testing only every 9-th character for whether it is `\q__str_recursion_tail` (also by converting 9 spaces at a time in the `\str_map_function:nN` case).

`\str_map_function:cN` For the `map_variable` functions we use a string assignment to store each character because spaces are made catcode 12 before the loop.

`\str_map_function:nN`

`\str_map_inline:NN`

`\str_map_inline:cn`

`\str_map_inline:nn`

`\str_map_variable:NNn`

`\str_map_variable:cNn`

`\str_map_variable:nNn`

`\str_map_break:`

`\str_map_break:n`

`__str_map_function:w`

`__str_map_function:Nn`

`__str_map_inline:NN`

`__str_map_variable:NNn`

```

4937 \cs_new:Npn \str_map_function:nN #1#2
4938 {
4939   \exp_after:wN \__str_map_function:w
4940   \exp_after:wN \__str_map_function:Nn \exp_after:wN #2
4941     \__kernel_tl_to_str:w {#1}
4942     \q__str_recursion_tail ? ~
4943   \prg_break_point:Nn \str_map_break: { }
4944 }
4945 \cs_new:Npn \str_map_function:NN
4946 { \exp_args:No \str_map_function:nN }
4947 \cs_new:Npn \__str_map_function:w #1 ~
4948 { #1 { ~ { ~ } \__str_map_function:w } }
4949 \cs_new:Npn \__str_map_function:Nn #1#2
4950 {
4951   \if_meaning:w \q__str_recursion_tail #2
4952     \exp_after:wN \str_map_break:
4953   \fi:
4954   #1 #2 \__str_map_function:Nn #1
4955 }
4956 \cs_generate_variant:Nn \str_map_function:NN { c }
4957 \cs_new_protected:Npn \str_map_inline:nn #1#2
4958 {
4959   \int_gincr:N \g__kernel_prg_map_int
4960   \cs_gset_protected:cpn
4961     { \__str_map_ \int_use:N \g__kernel_prg_map_int :w } ##1 {#2}
4962   \use:x

```



```

4963     {
4964         \exp_not:N \__str_map_inline:NN
4965         \exp_not:c { __str_map_ \int_use:N \g__kernel_prg_map_int :w }
4966         \__kernel_str_to_other_fast:n {#1}
4967     }
4968     \q__str_recursion_tail
4969     \prg_break_point:Nn \str_map_break:
4970     { \int_gdecr:N \g__kernel_prg_map_int }
4971 }
4972 \cs_new_protected:Npn \str_map_inline:Nn
4973 { \exp_args:No \str_map_inline:nn }
4974 \cs_generate_variant:Nn \str_map_inline:Nn { c }
4975 \cs_new:Npn \__str_map_inline:NN #1#2
4976 {
4977     \__str_if_recursion_tail_break:NN #2 \str_map_break:
4978     \exp_args:No #1 { \token_to_str:N #2 }
4979     \__str_map_inline:NN #1
4980 }
4981 \cs_new_protected:Npn \str_map_variable:nNn #1#2#3
4982 {
4983     \use:x
4984     {
4985         \exp_not:n { \__str_map_variable:NnN #2 {#3} }
4986         \__kernel_str_to_other_fast:n {#1}
4987     }
4988     \q__str_recursion_tail
4989     \prg_break_point:Nn \str_map_break: { }
4990 }
4991 \cs_new_protected:Npn \str_map_variable:NNn
4992 { \exp_args:No \str_map_variable:nNn }
4993 \cs_new_protected:Npn \__str_map_variable:NNn #1#2#3
4994 {
4995     \__str_if_recursion_tail_break:NN #3 \str_map_break:
4996     \str_set:Nn #1 {#3}
4997     \use:n {#2}
4998     \__str_map_variable:NnN #1 {#2}
4999 }
5000 \cs_generate_variant:Nn \str_map_variable:NNn { c }
5001 \cs_new:Npn \str_map_break:
5002 { \prg_map_break:Nn \str_map_break: { } }
5003 \cs_new:Npn \str_map_break:n
5004 { \prg_map_break:Nn \str_map_break: }

```

(End definition for `\str_map_function:NN` and others. These functions are documented on page 64.)

9.6 Accessing specific characters in a string

`__kernel_str_to_other:n` First apply `\tl_to_str:n`, then replace all spaces by “other” spaces, 8 at a time, storing the converted part of the string between the `\s__str_mark` and `\s__str_stop` markers. The end is detected when `__str_to_other_loop:w` finds one of the trailing A, distinguished from any contents of the initial token list by their category. Then `__str_to_other_end:w` is called, and finds the result between `\s__str_mark` and the first A (well, there is also the need to remove a space).

```

5005 \cs_new:Npn \__kernel_str_to_other:n #1
5006 {
5007   \exp_after:wN \__str_to_other_loop:w
5008   \tl_to_str:n {#1} ~ A ~ A ~ A ~ A ~ A ~ A ~ A ~ \s__str_mark \s__str_stop
5009 }
5010 \group_begin:
5011 \tex_lccode:D '\* = '\ %
5012 \tex_lccode:D '\A = '\A %
5013 \tex_lowercase:D
5014 {
5015   \group_end:
5016   \cs_new:Npn \__str_to_other_loop:w
5017     #1 ~ #2 ~ #3 ~ #4 ~ #5 ~ #6 ~ #7 ~ #8 ~ #9 \s__str_stop
5018   {
5019     \if_meaning:w A #8
5020     \__str_to_other_end:w
5021     \fi:
5022     \__str_to_other_loop:w
5023     #9 #1 * #2 * #3 * #4 * #5 * #6 * #7 * #8 * \s__str_stop
5024   }
5025   \cs_new:Npn \__str_to_other_end:w \fi: #1 \s__str_mark #2 * A #3 \s__str_stop
5026   { \fi: #2 }
5027 }

```

(End definition for __kernel_str_to_other:n, __str_to_other_loop:w, and __str_to_other_end:w.)

```

\__kernel_str_to_other_fast:n
\__kernel_str_to_other_fast_loop:w
\__str_to_other_fast_end:w

```

The difference with __kernel_str_to_other:n is that the converted part is left in the input stream, making these commands only restricted-expandable.

```

5028 \cs_new:Npn \__kernel_str_to_other_fast:n #1
5029 {
5030   \exp_after:wN \__str_to_other_fast_loop:w \tl_to_str:n {#1} ~
5031   A ~ A ~ A ~ A ~ A ~ A ~ A ~ A ~ A ~ \s__str_stop
5032 }
5033 \group_begin:
5034 \tex_lccode:D '\* = '\ %
5035 \tex_lccode:D '\A = '\A %
5036 \tex_lowercase:D
5037 {
5038   \group_end:
5039   \cs_new:Npn \__str_to_other_fast_loop:w
5040     #1 ~ #2 ~ #3 ~ #4 ~ #5 ~ #6 ~ #7 ~ #8 ~ #9 ~
5041   {
5042     \if_meaning:w A #9
5043     \__str_to_other_fast_end:w
5044     \fi:
5045     #1 * #2 * #3 * #4 * #5 * #6 * #7 * #8 * #9
5046     \__str_to_other_fast_loop:w *
5047   }
5048   \cs_new:Npn \__str_to_other_fast_end:w #1 * A #2 \s__str_stop {#1}
5049 }

```

(End definition for __kernel_str_to_other_fast:n, __kernel_str_to_other_fast_loop:w, and __str_to_other_fast_end:w.)

`\str_item:Nn` The `\str_item:nn` hands its argument with spaces escaped to `__str_item:nn`, and `\str_item:cn` makes sure to turn the result back into a proper string (with category code 10 spaces) eventually. The `\str_item_ignore_spaces:nn` function does not escape spaces, which are thus ignored by `__str_item:nn` since everything else is done with undelimited arguments. Evaluate the $\langle index \rangle$ argument #2 and count characters in the string, passing those two numbers to `__str_item:w` for further analysis. If the $\langle index \rangle$ is negative, shift it by the $\langle count \rangle$ to know the how many character to discard, and if that is still negative give an empty result. If the $\langle index \rangle$ is larger than the $\langle count \rangle$, give an empty result, and otherwise discard $\langle index \rangle - 1$ characters before returning the following one. The shift by -1 is obtained by inserting an empty brace group before the string in that case: that brace group also covers the case where the $\langle index \rangle$ is zero.

```

5050 \cs_new:Npn \str_item:Nn { \exp_args:No \str_item:nn }
5051 \cs_generate_variant:Nn \str_item:Nn { c }
5052 \cs_new:Npn \str_item:nn #1#2
5053 {
5054   \exp_args:Nf \tl_to_str:n
5055   {
5056     \exp_args:Nf \__str_item:nn
5057     { \__kernel_str_to_other:n {#1} } {#2}
5058   }
5059 }
5060 \cs_new:Npn \str_item_ignore_spaces:nn #1
5061 { \exp_args:No \__str_item:nn { \tl_to_str:n {#1} } }
5062 \cs_new:Npn \__str_item:nn #1#2
5063 {
5064   \exp_after:wN \__str_item:w
5065   \int_value:w \int_eval:n {#2} \exp_after:wN ;
5066   \int_value:w \__str_count:n {#1} ;
5067   #1 \s__str_stop
5068 }
5069 \cs_new:Npn \__str_item:w #1; #2;
5070 {
5071   \int_compare:nNnTF {#1} < 0
5072   {
5073     \int_compare:nNnTF {#1} < {-#2}
5074     { \__str_use_none_delimit_by_s_stop:w }
5075     {
5076       \exp_after:wN \__str_use_i_delimit_by_s_stop:nw
5077       \exp:w \exp_after:wN \__str_skip_exp_end:w
5078       \int_value:w \int_eval:n { #1 + #2 } ;
5079     }
5080   }
5081   {
5082     \int_compare:nNnTF {#1} > {#2}
5083     { \__str_use_none_delimit_by_s_stop:w }
5084     {
5085       \exp_after:wN \__str_use_i_delimit_by_s_stop:nw
5086       \exp:w \__str_skip_exp_end:w #1 ; { }
5087     }
5088   }
5089 }

```

(End definition for `\str_item:Nn` and others. These functions are documented on page 67.)

`__str_skip_exp_end:w` Removes `max(#1,0)` characters from the input stream, and then leaves `\exp_end:`. This should be expanded using `\exp:w`. We remove characters 8 at a time until there are at most 8 to remove. Then we do a dirty trick: the `\if_case:w` construction leaves between 0 and 8 times the `\or:` control sequence, and those `\or:` become arguments of `__str_skip_end:NNNNNNNN`. If the number of characters to remove is 6, say, then there are two `\or:` left, and the 8 arguments of `__str_skip_end:NNNNNNNN` are the two `\or:`, and 6 characters from the input stream, exactly what we wanted to remove. Then close the `\if_case:w` conditional with `\fi:`, and stop the initial expansion with `\exp_end:` (see places where `__str_skip_exp_end:w` is called).

```

5090 \cs_new:Npn \__str_skip_exp_end:w #1;
5091 {
5092   \if_int_compare:w #1 > 8 \exp_stop_f:
5093     \exp_after:wN \__str_skip_loop:wNNNNNNNN
5094   \else:
5095     \exp_after:wN \__str_skip_end:w
5096     \int_value:w \int_eval:w
5097   \fi:
5098   #1 ;
5099 }
5100 \cs_new:Npn \__str_skip_loop:wNNNNNNNN #1; #2#3#4#5#6#7#8#9
5101 {
5102   \exp_after:wN \__str_skip_exp_end:w
5103   \int_value:w \int_eval:n { #1 - 8 } ;
5104 }
5105 \cs_new:Npn \__str_skip_end:w #1 ;
5106 {
5107   \exp_after:wN \__str_skip_end:NNNNNNNN
5108   \if_case:w #1 \exp_stop_f: \or: \or: \or: \or: \or: \or: \or: \or:
5109 }
5110 \cs_new:Npn \__str_skip_end:NNNNNNNN #1#2#3#4#5#6#7#8 { \fi: \exp_end: }

```

(End definition for `__str_skip_exp_end:w` and others.)

`\str_range:Nnn` Sanitize the string. Then evaluate the arguments. At this stage we also decrement the `<start index>`, since our goal is to know how many characters should be removed. Then
`\str_range:nnn` limit the range to be non-negative and at most the length of the string (this avoids
`\str_range_ignore_spaces:nnn` needing to check for the end of the string when grabbing characters), shifting negative
`__str_range:nnn` numbers by the appropriate amount. Afterwards, skip characters, then keep some more,
`__str_range:w` and finally drop the end of the string.
`__str_range:nnw`

```

5111 \cs_new:Npn \str_range:Nnn { \exp_args:No \str_range:nnn }
5112 \cs_generate_variant:Nn \str_range:Nnn { c }
5113 \cs_new:Npn \str_range:nnn #1#2#3
5114 {
5115   \exp_args:Nf \tl_to_str:n
5116   {
5117     \exp_args:Nf \__str_range:nnn
5118     { \__kernel_str_to_other:n {#1} } {#2} {#3}
5119   }
5120 }
5121 \cs_new:Npn \str_range_ignore_spaces:nnn #1
5122 { \exp_args:No \__str_range:nnn { \tl_to_str:n {#1} } }
5123 \cs_new:Npn \__str_range:nnn #1#2#3
5124 {

```

```

5125     \exp_after:wN \_str_range:w
5126     \int_value:w \_str_count:n {#1} \exp_after:wN ;
5127     \int_value:w \int_eval:n { (#2) - 1 } \exp_after:wN ;
5128     \int_value:w \int_eval:n {#3} ;
5129     #1 \s__str_stop
5130 }
5131 \cs_new:Npn \_str_range:w #1; #2; #3;
5132 {
5133     \exp_args:Nf \_str_range:nnw
5134     { \_str_range_normalize:nn {#2} {#1} }
5135     { \_str_range_normalize:nn {#3} {#1} }
5136 }
5137 \cs_new:Npn \_str_range:nnw #1#2
5138 {
5139     \exp_after:wN \_str_collect_delimit_by_q_stop:w
5140     \int_value:w \int_eval:n { #2 - #1 } \exp_after:wN ;
5141     \exp:w \_str_skip_exp_end:w #1 ;
5142 }

```

(End definition for `\str_range:Nnn` and others. These functions are documented on page 68.)

`_str_range_normalize:nn` This function converts an *⟨index⟩* argument into an explicit position in the string (a result of 0 denoting “out of bounds”). Expects two explicit integer arguments: the *⟨index⟩* #1 and the string count #2. If #1 is negative, replace it by #1 + #2 + 1, then limit to the range [0, #2].

```

5143 \cs_new:Npn \_str_range_normalize:nn #1#2
5144 {
5145     \int_eval:n
5146     {
5147         \if_int_compare:w #1 < 0 \exp_stop_f:
5148         \if_int_compare:w #1 < -#2 \exp_stop_f:
5149         0
5150         \else:
5151         #1 + #2 + 1
5152         \fi:
5153     \else:
5154         \if_int_compare:w #1 < #2 \exp_stop_f:
5155         #1
5156         \else:
5157         #2
5158         \fi:
5159     \fi:
5160 }
5161 }

```

(End definition for `_str_range_normalize:nn`.)

`_str_collect_delimit_by_q_stop:w` Collects `max(#1,0)` characters, and removes everything else until `\s__str_stop`. This is somewhat similar to `_str_skip_exp_end:w`, but accepts integer expression arguments.
`_str_collect_loop:wn` This time we can only grab 7 characters at a time. At the end, we use an `\if_case:w` trick again, so that the 8 first arguments of `_str_collect_end:nnnnnnnnnw` are some `\or:`, followed by an `\fi:`, followed by #1 characters from the input stream. Simply leaving this in the input stream closes the conditional properly and the `\or:` disappear.
`_str_collect_loop:wnnnnnnnnw`
`_str_collect_end:wn`
`_str_collect_end:nnnnnnnnnw`

```

5162 \cs_new:Npn \_str_collect_delimit_by_q_stop:w #1;

```

```

5163 { \__str_collect_loop:wn #1 ; { } }
5164 \cs_new:Npn \__str_collect_loop:wn #1 ;
5165 {
5166   \if_int_compare:w #1 > 7 \exp_stop_f:
5167   \exp_after:wN \__str_collect_loop:wnNNNNNNN
5168   \else:
5169   \exp_after:wN \__str_collect_end:wn
5170   \fi:
5171   #1 ;
5172 }
5173 \cs_new:Npn \__str_collect_loop:wnNNNNNNN #1; #2 #3#4#5#6#7#8#9
5174 {
5175   \exp_after:wN \__str_collect_loop:wn
5176   \int_value:w \int_eval:n { #1 - 7 } ;
5177   { #2 #3#4#5#6#7#8#9 }
5178 }
5179 \cs_new:Npn \__str_collect_end:wn #1 ;
5180 {
5181   \exp_after:wN \__str_collect_end:nnnnnnnnw
5182   \if_case:w \if_int_compare:w #1 > 0 \exp_stop_f:
5183   #1 \else: 0 \fi: \exp_stop_f:
5184   \or: \or: \or: \or: \or: \or: \or: \fi:
5185 }
5186 \cs_new:Npn \__str_collect_end:nnnnnnnnw #1#2#3#4#5#6#7#8 #9 \s__str_stop
5187 { #1#2#3#4#5#6#7#8 }

```

(End definition for __str_collect_delimit_by_q_stop:w and others.)

9.7 Counting characters

\str_count_spaces:N To speed up this function, we grab and discard 9 space-delimited arguments in each iteration of the loop. The loop stops when the last argument is one of the trailing $X\langle number \rangle$, and that $\langle number \rangle$ is added to the sum of 9 that precedes, to adjust the result.

```

\__str_count_spaces_loop:w
5188 \cs_new:Npn \str_count_spaces:N
5189 { \exp_args:No \str_count_spaces:n }
5190 \cs_generate_variant:Nn \str_count_spaces:N { c }
5191 \cs_new:Npn \str_count_spaces:n #1
5192 {
5193   \int_eval:n
5194   {
5195     \exp_after:wN \__str_count_spaces_loop:w
5196     \tl_to_str:n {#1} ~
5197     X 7 ~ X 6 ~ X 5 ~ X 4 ~ X 3 ~ X 2 ~ X 1 ~ X 0 ~ X -1 ~
5198     \s__str_stop
5199   }
5200 }
5201 \cs_new:Npn \__str_count_spaces_loop:w #1~#2~#3~#4~#5~#6~#7~#8~#9~
5202 {
5203   \if_meaning:w X #9
5204   \__str_use_i_delimit_by_s_stop:nw
5205   \fi:
5206   9 + \__str_count_spaces_loop:w
5207 }

```

(End definition for `\str_count_spaces:N`, `\str_count_spaces:n`, and `__str_count_spaces_loop:w`. These functions are documented on page 66.)

`\str_count:N` To count characters in a string we could first escape all spaces using `__kernel_str_to_other:n`, then pass the result to `\tl_count:n`. However, the escaping step would be quadratic in the number of characters in the string, and we can do better. Namely, sum the number of spaces (`\str_count_spaces:n`) and the result of `\tl_count:n`, which ignores spaces. Since strings tend to be longer than token lists, we use specialized functions to count characters ignoring spaces. Namely, `loop`, grabbing 9 non-space characters at each step, and end as soon as we reach one of the 9 trailing items. The internal function `__str_count:n`, used in `\str_item:nn` and `\str_range:nnn`, is similar to `\str_count_ignore_spaces:n` but expects its argument to already be a string or a string with spaces escaped.

```

5208 \cs_new:Npn \str_count:N { \exp_args:No \str_count:n }
5209 \cs_generate_variant:Nn \str_count:N { c }
5210 \cs_new:Npn \str_count:n #1
5211 {
5212   \__str_count_aux:n
5213   {
5214     \str_count_spaces:n {#1}
5215     + \exp_after:wN \__str_count_loop:NNNNNNNNN \tl_to_str:n {#1}
5216   }
5217 }
5218 \cs_new:Npn \__str_count:n #1
5219 {
5220   \__str_count_aux:n
5221   { \__str_count_loop:NNNNNNNNN #1 }
5222 }
5223 \cs_new:Npn \str_count_ignore_spaces:n #1
5224 {
5225   \__str_count_aux:n
5226   { \exp_after:wN \__str_count_loop:NNNNNNNNN \tl_to_str:n {#1} }
5227 }
5228 \cs_new:Npn \__str_count_aux:n #1
5229 {
5230   \int_eval:n
5231   {
5232     #1
5233     { X 8 } { X 7 } { X 6 }
5234     { X 5 } { X 4 } { X 3 }
5235     { X 2 } { X 1 } { X 0 }
5236     \s__str_stop
5237   }
5238 }
5239 \cs_new:Npn \__str_count_loop:NNNNNNNNN #1#2#3#4#5#6#7#8#9
5240 {
5241   \if_meaning:w X #9
5242     \exp_after:wN \__str_use_none_delimit_by_s_stop:w
5243   \fi:
5244   9 + \__str_count_loop:NNNNNNNNN
5245 }
```

(End definition for `\str_count:N` and others. These functions are documented on page 66.)

9.8 The first character in a string

`\str_head:N` The `_ignore_spaces` variant applies `\tl_to_str:n` then grabs the first item, thus skipping spaces. As usual, `\str_head:N` expands its argument and hands it to `\str_head:n`.
`\str_head:c` To circumvent the fact that TeX skips spaces when grabbing undelimited macro parameters, `__str_head:w` takes an argument delimited by a space. If #1 starts with a non-space character, `__str_use_i_delimit_by_s_stop:nw` leaves that in the input stream. On the other hand, if #1 starts with a space, the `__str_head:w` takes an empty argument, and the single (initially braced) space in the definition of `__str_head:w` makes its way to the output. Finally, for an empty argument, the (braced) empty brace group in the definition of `\str_head:n` gives an empty result after passing through `__str_use_i_delimit_by_s_stop:nw`.
`\str_head:n`

```

5246 \cs_new:Npn \str_head:N { \exp_args:No \str_head:n }
5247 \cs_generate_variant:Nn \str_head:N { c }
5248 \cs_new:Npn \str_head:n #1
5249 {
5250   \exp_after:wN __str_head:w
5251   \tl_to_str:n {#1}
5252   { { } } ~ \s__str_stop
5253 }
5254 \cs_new:Npn __str_head:w #1 ~ %
5255 { \__str_use_i_delimit_by_s_stop:nw #1 { ~ } }
5256 \cs_new:Npn \str_head_ignore_spaces:n #1
5257 {
5258   \exp_after:wN __str_use_i_delimit_by_s_stop:nw
5259   \tl_to_str:n {#1} { } \s__str_stop
5260 }
```

(End definition for `\str_head:N` and others. These functions are documented on page 67.)

`\str_tail:N` Getting the tail is a little bit more convoluted than the head of a string. We hit the front of the string with `\reverse_if:N` `\if_charcode:w` `\scan_stop:.` This removes the first character, and necessarily makes the test true, since the character cannot match `\scan_stop:.` The auxiliary function then inserts the required `\fi:` to close the conditional, and leaves the tail of the string in the input stream. The details are such that an empty string has an empty tail (this requires in particular that the end-marker X be unexpandable and not a control sequence). The `_ignore_spaces` is rather simpler: after converting the input to a string, `__str_tail_auxii:w` removes one undelimited argument and leaves everything else until an end-marker `\s__str_mark`. One can check that an empty (or blank) string yields an empty tail.
`\str_tail:c`
`\str_tail:n`
`\str_tail_ignore_spaces:n`
`__str_tail_auxi:w`
`__str_tail_auxii:w`

```

5261 \cs_new:Npn \str_tail:N { \exp_args:No \str_tail:n }
5262 \cs_generate_variant:Nn \str_tail:N { c }
5263 \cs_new:Npn \str_tail:n #1
5264 {
5265   \exp_after:wN __str_tail_auxi:w
5266   \reverse_if:N \if_charcode:w
5267   \scan_stop: \tl_to_str:n {#1} X X \s__str_stop
5268 }
5269 \cs_new:Npn __str_tail_auxi:w #1 X #2 \s__str_stop { \fi: #1 }
5270 \cs_new:Npn \str_tail_ignore_spaces:n #1
5271 {
5272   \exp_after:wN __str_tail_auxii:w
5273   \tl_to_str:n {#1} \s__str_mark \s__str_mark \s__str_stop
```



```

5274 }
5275 \cs_new:Npn \__str_tail_auxii:w #1 #2 \s__str_mark #3 \s__str_stop { #2 }

```

(End definition for `\str_tail:N` and others. These functions are documented on page 67.)

9.9 String manipulation

```

\str_foldcase:n Case changing for programmatic reasons is done by first detokenizing input then doing
\str_foldcase:V a simple loop that only has to worry about spaces and everything else. The output is
\str_lowercase:n detokenized to allow data sharing with text-based case changing.
\str_lowercase:f
\str_uppercase:n
\str_uppercase:f
\__str_change_case:nn
\__str_change_case_aux:nn
\__str_change_case_result:n
\__str_change_case_output:nw
\__str_change_case_output:fw
\__str_change_case_end:nw
\__str_change_case_loop:nw
\__str_change_case_space:n
\__str_change_case_char:nN
5276 \cs_new:Npn \str_foldcase:n #1 { \__str_change_case:nn {#1} { fold } }
5277 \cs_new:Npn \str_lowercase:n #1 { \__str_change_case:nn {#1} { lower } }
5278 \cs_new:Npn \str_uppercase:n #1 { \__str_change_case:nn {#1} { upper } }
5279 \cs_generate_variant:Nn \str_foldcase:n { V }
5280 \cs_generate_variant:Nn \str_lowercase:n { f }
5281 \cs_generate_variant:Nn \str_uppercase:n { f }
5282 \cs_new:Npn \__str_change_case:nn #1
5283 {
5284   \exp_after:wN \__str_change_case_aux:nn \exp_after:wN
5285   { \tl_to_str:n {#1} }
5286 }
5287 \cs_new:Npn \__str_change_case_aux:nn #1#2
5288 {
5289   \__str_change_case_loop:nw {#2} #1 \q__str_recursion_tail \q__str_recursion_stop
5290   \__str_change_case_result:n { }
5291 }
5292 \cs_new:Npn \__str_change_case_output:nw #1#2 \__str_change_case_result:n #3
5293 { #2 \__str_change_case_result:n { #3 #1 } }
5294 \cs_generate_variant:Nn \__str_change_case_output:nw { f }
5295 \cs_new:Npn \__str_change_case_end:wn #1 \__str_change_case_result:n #2
5296 { \tl_to_str:n {#2} }
5297 \cs_new:Npn \__str_change_case_loop:nw #1#2 \q__str_recursion_stop
5298 {
5299   \tl_if_head_is_space:nTF {#2}
5300   { \__str_change_case_space:n }
5301   { \__str_change_case_char:nN }
5302   {#1} #2 \q__str_recursion_stop
5303 }
5304 \exp_last_unbraced:NNNNo
5305 \cs_new:Npn \__str_change_case_space:n #1 \c_space_tl
5306 {
5307   \__str_change_case_output:nw { ~ }
5308   \__str_change_case_loop:nw {#1}
5309 }
5310 \cs_new:Npn \__str_change_case_char:nN #1#2
5311 {
5312   \__str_if_recursion_tail_stop_do:Nn #2
5313   { \__str_change_case_end:wn }
5314   \__str_change_case_output:fw
5315   { \use:c { char_str_ #1 case:N } #2 }
5316   \__str_change_case_loop:nw {#1}
5317 }

```

(End definition for `\str_foldcase:n` and others. These functions are documented on page 70.)

`\c_ampersand_str` For all of those strings, use `\cs_to_str:N` to get characters with the correct category code without worries

`\c_atsign_str`

`\c_backslash_str` 5318 `\str_const:Nx \c_ampersand_str { \cs_to_str:N \& }`

`\c_left_brace_str` 5319 `\str_const:Nx \c_atsign_str { \cs_to_str:N \@ }`

`\c_right_brace_str` 5320 `\str_const:Nx \c_backslash_str { \cs_to_str:N \\ }`

`\c_circumflex_str` 5321 `\str_const:Nx \c_left_brace_str { \cs_to_str:N \{ }`

`\c_colon_str` 5322 `\str_const:Nx \c_right_brace_str { \cs_to_str:N \} }`

`\c_dollar_str` 5323 `\str_const:Nx \c_circumflex_str { \cs_to_str:N \^ }`

`\c_hash_str` 5324 `\str_const:Nx \c_colon_str { \cs_to_str:N \: }`

`\c_percent_str` 5325 `\str_const:Nx \c_dollar_str { \cs_to_str:N \$ }`

`\c_tilde_str` 5326 `\str_const:Nx \c_hash_str { \cs_to_str:N \# }`

`\c_underscore_str` 5327 `\str_const:Nx \c_percent_str { \cs_to_str:N \% }`

`\c_zero_str` 5328 `\str_const:Nx \c_tilde_str { \cs_to_str:N \~ }`

5329 `\str_const:Nx \c_underscore_str { \cs_to_str:N _ }`

5330 `\str_const:Nx \c_zero_str { 0 }`

(End definition for `\c_ampersand_str` and others. These variables are documented on page 71.)

`\l_tmpa_str` Scratch strings.

`\l_tmpb_str` 5331 `\str_new:N \l_tmpa_str`

`\g_tmpa_str` 5332 `\str_new:N \l_tmpb_str`

`\g_tmpb_str` 5333 `\str_new:N \g_tmpa_str`

5334 `\str_new:N \g_tmpb_str`

(End definition for `\l_tmpa_str` and others. These variables are documented on page 71.)

9.10 Viewing strings

`\str_show:n` Displays a string on the terminal.

`\str_show:N` 5335 `\cs_new_eq:NN \str_show:n \tl_show:n`

`\str_show:c` 5336 `\cs_new_eq:NN \str_show:N \tl_show:N`

`\str_log:n` 5337 `\cs_generate_variant:Nn \str_show:N { c }`

`\str_log:N` 5338 `\cs_new_eq:NN \str_log:n \tl_log:n`

`\str_log:c` 5339 `\cs_new_eq:NN \str_log:N \tl_log:N`

5340 `\cs_generate_variant:Nn \str_log:N { c }`

(End definition for `\str_show:n` and others. These functions are documented on page 70.)

5341 `\</package>`

10 l3str-convert implementation

5342 `*package>`

5343 `\@@=str>`

10.1 Helpers

10.1.1 Variables and constants

`__str_tmp:w` Internal scratch space for some functions.

`\l__str_internal_tl` 5344 `\cs_new_protected:Npn __str_tmp:w { }`

5345 `\tl_new:N \l__str_internal_tl`

(End definition for `__str_tmp:w` and `\l__str_internal_tl`.)

`\g__str_result_tl` The `\g__str_result_tl` variable is used to hold the result of various internal string operations (mostly conversions) which are typically performed in a group. The variable is global so that it remains defined outside the group, to be assigned to a user-provided variable.

```
5346 \tl_new:N \g__str_result_tl
```

(End definition for `\g__str_result_tl`.)

`\c__str_replacement_char_int` When converting, invalid bytes are replaced by the Unicode replacement character "FFFD.

```
5347 \int_const:Nn \c__str_replacement_char_int { "FFFD }
```

(End definition for `\c__str_replacement_char_int`.)

`\c__str_max_byte_int` The maximal byte number.

```
5348 \int_const:Nn \c__str_max_byte_int { 255 }
```

(End definition for `\c__str_max_byte_int`.)

`\s__str` Internal scan marks.

```
5349 \scan_new:N \s__str
```

(End definition for `\s__str`.)

`\q__str_nil` Internal quarks.

```
5350 \quark_new:N \q__str_nil
```

(End definition for `\q__str_nil`.)

`\g__str_alias_prop` To avoid needing one file per encoding/escaping alias, we keep track of those in a property list.

```
5351 \prop_new:N \g__str_alias_prop
5352 \prop_gput:Nnn \g__str_alias_prop { latin1 } { iso88591 }
5353 \prop_gput:Nnn \g__str_alias_prop { latin2 } { iso88592 }
5354 \prop_gput:Nnn \g__str_alias_prop { latin3 } { iso88593 }
5355 \prop_gput:Nnn \g__str_alias_prop { latin4 } { iso88594 }
5356 \prop_gput:Nnn \g__str_alias_prop { latin5 } { iso88599 }
5357 \prop_gput:Nnn \g__str_alias_prop { latin6 } { iso885910 }
5358 \prop_gput:Nnn \g__str_alias_prop { latin7 } { iso885913 }
5359 \prop_gput:Nnn \g__str_alias_prop { latin8 } { iso885914 }
5360 \prop_gput:Nnn \g__str_alias_prop { latin9 } { iso885915 }
5361 \prop_gput:Nnn \g__str_alias_prop { latin10 } { iso885916 }
5362 \prop_gput:Nnn \g__str_alias_prop { utf16le } { utf16 }
5363 \prop_gput:Nnn \g__str_alias_prop { utf16be } { utf16 }
5364 \prop_gput:Nnn \g__str_alias_prop { utf32le } { utf32 }
5365 \prop_gput:Nnn \g__str_alias_prop { utf32be } { utf32 }
5366 \prop_gput:Nnn \g__str_alias_prop { hexadecimal } { hex }
```

(End definition for `\g__str_alias_prop`.)

`\g__str_error_bool` In conversion functions with a built-in conditional, errors are not reported directly to the user, but the information is collected in this boolean, used at the end to decide on which branch of the conditional to take.

```
5367 \bool_new:N \g__str_error_bool
```

(End definition for `\g__str_error_bool`.)

str_byte Conversions from one *<encoding>/<escaping>* pair to another are done within x-expanding assignments. Errors are signalled by raising the relevant flag.

```
5368 \flag_new:n { str_byte }
5369 \flag_new:n { str_error }
```

(End definition for `str_byte` and `str_error`. These variables are documented on page ??.)

10.2 String conditionals

```
\__str_if_contains_char:NnT      \__str_if_contains_char:nnTF {<token list>} <char>
\__str_if_contains_char:NnTF      Expects the <token list> to be an <other string>: the caller is responsible for ensuring
\__str_if_contains_char:nnTF      that no (too-)special catcodes remain. Loop over the characters of the string, comparing
    \__str_if_contains_char_aux:nn character codes. The loop is broken if character codes match. Otherwise we return
    \__str_if_contains_char_auxi:nN “false”.
    \__str_if_contains_char_true:

5370 \prg_new_conditional:Npnn \__str_if_contains_char:Nn #1#2 { T , TF }
5371 {
5372     \exp_after:wN \__str_if_contains_char_aux:nn \exp_after:wN {#1} {#2}
5373     { \prg_break:n { ? \fi: } }
5374     \prg_break_point:
5375     \prg_return_false:
5376 }
5377 \cs_new:Npn \__str_if_contains_char_aux:nn #1#2
5378 { \__str_if_contains_char_auxi:nN {#2} #1 }
5379 \prg_new_conditional:Npnn \__str_if_contains_char:nn #1#2 { TF }
5380 {
5381     \__str_if_contains_char_auxi:nN {#2} #1 { \prg_break:n { ? \fi: } }
5382     \prg_break_point:
5383     \prg_return_false:
5384 }
5385 \cs_new:Npn \__str_if_contains_char_auxi:nN #1#2
5386 {
5387     \if_charcode:w #1 #2
5388     \exp_after:wN \__str_if_contains_char_true:
5389     \fi:
5390     \__str_if_contains_char_auxi:nN {#1}
5391 }
5392 \cs_new:Npn \__str_if_contains_char_true:
5393 { \prg_break:n { \prg_return_true: \use_none:n } }
```

(End definition for `__str_if_contains_char:NnT` and others.)

__str_octal_use:NNTF `__str_octal_use:NNTF <token> {<true code>} {<false code>}`
 If the *<token>* is an octal digit, it is left in the input stream, *followed* by the *<true code>*. Otherwise, the *<false code>* is left in the input stream.

T_EXhackers note: This function will fail if the escape character is an octal digit. We are thus careful to set the escape character to a known value before using it. T_EX dutifully detects

octal digits for us: if #1 is an octal digit, then the right-hand side of the comparison is '1#1, greater than 1. Otherwise, the right-hand side stops as '1, and the conditional takes the false branch.

```

5394 \prg_new_conditional:Npnn \__str_octal_use:N #1 { TF }
5395 {
5396   \if_int_compare:w 1 < '1 \token_to_str:N #1 \exp_stop_f:
5397     #1 \prg_return_true:
5398   \else:
5399     \prg_return_false:
5400   \fi:
5401 }

```

(End definition for __str_octal_use:NTF.)

__str_hexadecimal_use:NTF TeX detects uppercase hexadecimal digits for us (see __str_octal_use:NTF), but not the lowercase letters, which we need to detect and replace by their uppercase counterpart.

```

5402 \prg_new_conditional:Npnn \__str_hexadecimal_use:N #1 { TF }
5403 {
5404   \if_int_compare:w 1 < "1 \token_to_str:N #1 \exp_stop_f:
5405     #1 \prg_return_true:
5406   \else:
5407     \if_case:w \int_eval:n { \exp_after:wN ' \token_to_str:N #1 - 'a }
5408       A
5409     \or: B
5410     \or: C
5411     \or: D
5412     \or: E
5413     \or: F
5414     \else:
5415       \prg_return_false:
5416       \exp_after:wN \use_none:n
5417     \fi:
5418     \prg_return_true:
5419   \fi:
5420 }

```

(End definition for __str_hexadecimal_use:NTF.)

10.3 Conversions

10.3.1 Producing one byte or character

\c__str_byte_0_tl For each integer N in the range $[0, 255]$, we create a constant token list which holds three character tokens with category code other: the character with character code N , followed by the representation of N as two hexadecimal digits. The value -1 is given a default token list which ensures that later functions give an empty result for the input -1 .

```

5421 \group_begin:
5422   \__kernel_tl_set:Nx \l__str_internal_tl { \tl_to_str:n { 0123456789ABCDEF } }
5423   \tl_map_inline:Nn \l__str_internal_tl
5424     {
5425       \tl_map_inline:Nn \l__str_internal_tl
5426         {
5427           \tl_const:cx { c__str_byte_ \int_eval:n {"#1##1} _tl }
5428             { \char_generate:nn { "#1##1 } { 12 } #1 ##1 }
5429         }
5430     }
5431 \group_end:
5432 \tl_const:cn { c__str_byte_-1_tl } { { } \use_none:n { } }

```

(End definition for `\c__str_byte_0_t1` and others.)

`__str_output_byte:n` Those functions must be used carefully: feeding them a value outside the range $[-1, 255]$
`__str_output_byte:w` will attempt to use the undefined token list variable `\c__str_byte_⟨number⟩_t1`. As-
`__str_output_hexadecimal:n` summing that the argument is in the right range, we expand the corresponding token list,
`__str_output_end:` and pick either the byte (first token) or the hexadecimal representations (second and
third tokens). The value -1 produces an empty result in both cases.

```

5433 \cs_new:Npn \__str_output_byte:n #1
5434   { \__str_output_byte:w #1 \__str_output_end: }
5435 \cs_new:Npn \__str_output_byte:w
5436   {
5437     \exp_after:wN \exp_after:wN
5438     \exp_after:wN \use_i:nnn
5439     \cs:w c__str_byte_ \int_eval:w
5440   }
5441 \cs_new:Npn \__str_output_hexadecimal:n #1
5442   {
5443     \exp_after:wN \exp_after:wN
5444     \exp_after:wN \use_none:n
5445     \cs:w c__str_byte_ \int_eval:n {#1} _t1 \cs_end:
5446   }
5447 \cs_new:Npn \__str_output_end:
5448   { \scan_stop: _t1 \cs_end: }

```

(End definition for `__str_output_byte:n` and others.)

`__str_output_byte_pair_be:n` Convert a number in the range $[0, 65535]$ to a pair of bytes, either big-endian or little-
`__str_output_byte_pair_le:n` endian.
`__str_output_byte_pair:nnN`

```

5449 \cs_new:Npn \__str_output_byte_pair_be:n #1
5450   {
5451     \exp_args:Nf \__str_output_byte_pair:nnN
5452     { \int_div_truncate:nn { #1 } { "100 } } {#1} \use:nn
5453   }
5454 \cs_new:Npn \__str_output_byte_pair_le:n #1
5455   {
5456     \exp_args:Nf \__str_output_byte_pair:nnN
5457     { \int_div_truncate:nn { #1 } { "100 } } {#1} \use_ii_i:nn
5458   }
5459 \cs_new:Npn \__str_output_byte_pair:nnN #1#2#3
5460   {
5461     #3
5462     { \__str_output_byte:n { #1 } }
5463     { \__str_output_byte:n { #2 - #1 * "100 } }
5464   }

```

(End definition for `__str_output_byte_pair_be:n`, `__str_output_byte_pair_le:n`, and `__str_output_byte_pair:nnN`.)

10.3.2 Mapping functions for conversions

`__str_convert_gmap:N` This maps the function `#1` over all characters in `\g__str_result_t1`, which should be a
`__str_convert_gmap_loop:NN` byte string in most cases, sometimes a native string.

```

5465 \cs_new_protected:Npn \__str_convert_gmap:N #1
5466   {

```

```

5467     \__kernel_tl_gset:Nx \g__str_result_tl
5468     {
5469         \exp_after:wN \__str_convert_gmap_loop:NN
5470         \exp_after:wN #1
5471         \g__str_result_tl { ? \prg_break: }
5472         \prg_break_point:
5473     }
5474 }
5475 \cs_new:Npn \__str_convert_gmap_loop:NN #1#2
5476 {
5477     \use_none:n #2
5478     #1#2
5479     \__str_convert_gmap_loop:NN #1
5480 }

```

(End definition for __str_convert_gmap:N and __str_convert_gmap_loop:NN.)

__str_convert_gmap_internal:N
__str_convert_gmap_internal_loop:Nw

This maps the function #1 over all character codes in \g__str_result_tl, which must be in the internal representation.

```

5481 \cs_new_protected:Npn \__str_convert_gmap_internal:N #1
5482 {
5483     \__kernel_tl_gset:Nx \g__str_result_tl
5484     {
5485         \exp_after:wN \__str_convert_gmap_internal_loop:Nww
5486         \exp_after:wN #1
5487         \g__str_result_tl \s__str \s__str_stop \prg_break: \s__str
5488         \prg_break_point:
5489     }
5490 }
5491 \cs_new:Npn \__str_convert_gmap_internal_loop:Nww #1 #2 \s__str #3 \s__str
5492 {
5493     \__str_use_none_delimit_by_s_stop:w #3 \s__str_stop
5494     #1 {#3}
5495     \__str_convert_gmap_internal_loop:Nww #1
5496 }

```

(End definition for __str_convert_gmap_internal:N and __str_convert_gmap_internal_loop:Nw.)

10.3.3 Error-reporting during conversion

__str_if_flag_error:nnx
__str_if_flag_no_error:nnx

When converting using the function \str_set_convert:Nnnn, errors should be reported to the user after each step in the conversion. Errors are signalled by raising some flag (typically @@_error), so here we test that flag: if it is raised, give the user an error, otherwise remove the arguments. On the other hand, in the conditional functions \str_set_convert:NnnnTF, errors should be suppressed. This is done by changing __str_if_flag_error:nnx into __str_if_flag_no_error:nnx locally.

```

5497 \cs_new_protected:Npn \__str_if_flag_error:nnx #1
5498 {
5499     \flag_if_raised:nTF {#1}
5500     { \__kernel_msg_error:nnx { str } }
5501     { \use_none:nn }
5502 }
5503 \cs_new_protected:Npn \__str_if_flag_no_error:nnx #1#2#3
5504 { \flag_if_raised:nT {#1} { \bool_gset_true:N \g__str_error_bool } }

```

(End definition for `_str_if_flag_error:nxx` and `_str_if_flag_no_error:nxx`.)

`_str_if_flag_times:nT` At the end of each conversion step, we raise all relevant errors as one error message, built on the fly. The height of each flag indicates how many times a given error was encountered. This function prints `#2` followed by the number of occurrences of an error if it occurred, nothing otherwise.

```
5505 \cs_new:Npn \_str_if_flag_times:nT #1#2
5506   { \flag_if_raised:nT {#1} { #2~(x \flag_height:n {#1} ) } }
```

(End definition for `_str_if_flag_times:nT`.)

10.3.4 Framework for conversions

Most functions in this module expect to be working with “native” strings. Strings can also be stored as bytes, in one of many encodings, for instance UTF8. The bytes themselves can be expressed in various ways in terms of \TeX tokens, for instance as pairs of hexadecimal digits. The questions of going from arbitrary Unicode code points to bytes, and from bytes to tokens are mostly independent.

Conversions are done in four steps:

- “unescape” produces a string of bytes;
- “decode” takes in a string of bytes, and converts it to a list of Unicode characters in an internal representation, with items of the form

$\langle bytes \rangle \backslash s_str \langle Unicode\ code\ point \rangle \backslash s_str$

where we have collected the $\langle bytes \rangle$ which combined to form this particular Unicode character, and the $\langle Unicode\ code\ point \rangle$ is in the range $[0, "10FFFF]$.

- “encode” encodes the internal list of code points as a byte string in the new encoding;
- “escape” escapes bytes as requested.

The process is modified in case one of the encoding is empty (or the conversion function has been set equal to the empty encoding because it was not found): then the unescape or escape step is ignored, and the decode or encode steps work on tokens instead of bytes. Otherwise, each step must ensure that it passes a correct byte string or internal string to the next step.

`\str_set_convert:Nnnn` The input string is stored in `\g__str_result_tl`, then we: unescape and decode; encode and escape; exit the group and store the result in the user’s variable. The various conversion functions all act on `\g__str_result_tl`. Errors are silenced for the conditional functions by redefining `_str_if_flag_error:nxx` locally.

```
\str_gset_convert:Nnnn
\str_set_convert:NnnnTF
\str_gset_convert:NnnnTF
\_str_convert:nNNnnn
5507 \cs_new_protected:Npn \str_set_convert:Nnnn
5508   { \_str_convert:nNNnnn { } \tl_set_eq:NN }
5509 \cs_new_protected:Npn \str_gset_convert:Nnnn
5510   { \_str_convert:nNNnnn { } \tl_gset_eq:NN }
5511 \prg_new_protected_conditional:Npnn
5512   \str_set_convert:Nnnn #1#2#3#4 { T , F , TF }
5513   {
5514     \bool_gset_false:N \g__str_error_bool
5515     \_str_convert:nNNnnn
```



```

5516     { \cs_set_eq:NN \__str_if_flag_error:nxx \__str_if_flag_no_error:nxx }
5517     \tl_set_eq:NN #1 {#2} {#3} {#4}
5518     \bool_if:NTF \g__str_error_bool \prg_return_false: \prg_return_true:
5519   }
5520 \prg_new_protected_conditional:Npnn
5521   \str_gset_convert:Nnnn #1#2#3#4 { T , F , TF }
5522   {
5523     \bool_gset_false:N \g__str_error_bool
5524     \__str_convert:nNNnnn
5525     { \cs_set_eq:NN \__str_if_flag_error:nxx \__str_if_flag_no_error:nxx }
5526     \tl_gset_eq:NN #1 {#2} {#3} {#4}
5527     \bool_if:NTF \g__str_error_bool \prg_return_false: \prg_return_true:
5528   }
5529 \cs_new_protected:Npn \__str_convert:nNNnnn #1#2#3#4#5#6
5530   {
5531     \group_begin:
5532     #1
5533     \__kernel_tl_gset:Nx \g__str_result_tl { \__kernel_str_to_other_fast:n {#4} }
5534     \exp_after:wN \__str_convert:wwwnn
5535     \tl_to_str:n {#5} /// \s__str_stop
5536     { decode } { unescape }
5537     \prg_do_nothing:
5538     \__str_convert_decode_:
5539     \exp_after:wN \__str_convert:wwwnn
5540     \tl_to_str:n {#6} /// \s__str_stop
5541     { encode } { escape }
5542     \use_ii_i:nn
5543     \__str_convert_encode_:
5544     \group_end:
5545     #2 #3 \g__str_result_tl
5546   }

```

(End definition for `\str_set_convert:Nnnn` and others. These functions are documented on page 72.)

`__str_convert:wwwnn` The task of `__str_convert:wwwnn` is to split *⟨encoding⟩/⟨escaping⟩* pairs into their components, #1 and #2. Calls to `__str_convert:nnn` ensure that the corresponding conversion functions are defined. The third auxiliary does the main work.

- #1 is the encoding conversion function;
- #2 is the escaping function;
- #3 is the escaping name for use in an error message;
- #4 is `\prg_do_nothing:` for unescaping/decoding, and `\use_ii_i:nn` for encoding/escaping;
- #5 is the default encoding function (either “decode” or “encode”), for which there should be no escaping.

Let us ignore the native encoding for a second. In the unescaping/decoding phase, we want to do #2#1 in this order, and in the encoding/escaping phase, the order should be reversed: #4#2#1 does exactly that. If one of the encodings is the default (native), then the escaping should be ignored, with an error if any was given, and only the encoding, #1, should be performed.

```

5547 \cs_new_protected:Npn \__str_convert:wwwnn
5548   #1 / #2 // #3 \s__str_stop #4#5
5549   {
5550     \__str_convert:nnn {enc} {#4} {#1}
5551     \__str_convert:nnn {esc} {#5} {#2}
5552     \exp_args:Ncc \__str_convert:NNnNN
5553       { __str_convert_#4_#1: } { __str_convert_#5_#2: } {#2}
5554   }
5555 \cs_new_protected:Npn \__str_convert:NNnNN #1#2#3#4#5
5556   {
5557     \if_meaning:w #1 #5
5558       \tl_if_empty:nF {#3}
5559       { \__kernel_msg_error:nxx { str } { native-escaping } {#3} }
5560       #1
5561     \else:
5562       #4 #2 #1
5563     \fi:
5564   }

```

(End definition for __str_convert:wwwnn and __str_convert:NNnNN.)

__str_convert:nnn The arguments of __str_convert:nnn are: `enc` or `esc`, used to build filenames, the type of the conversion (unescape, decode, encode, escape), and the encoding or escaping name. If the function is already defined, no need to do anything. Otherwise, filter out all non-alphanumerics in the name, and lowercase it. Feed that, and the same three arguments, to __str_convert:nnnn. The task is then to make sure that the conversion function `#3_#1` corresponding to the type `#3` and filtered name `#1` is defined, then set our initial conversion function `#3_#4` equal to that.

How do we get the `#3_#1` conversion to be defined if it isn't? Two main cases.

First, if `#1` is a key in `\g__str_alias_prop`, then the value `\l__str_internal_tl` tells us what file to load. Loading is skipped if the file was already read, *i.e.*, if the conversion command based on `\l__str_internal_tl` already exists. Otherwise, try to load the file; if that fails, there is an error, use the default empty name instead.

Second, `#1` may be absent from the property list. The `\cs_if_exist:cF` test is automatically false, and we search for a file defining the encoding or escaping `#1` (this should allow third-party `.def` files). If the file is not found, there is an error, use the default empty name instead.

In all cases, the conversion based on `\l__str_internal_tl` is defined, so we can set the `#3_#1` function equal to that. In some cases (*e.g.*, `utf16be`), the `#3_#1` function is actually defined within the file we just loaded, and it is different from the `\l__str_internal_tl`-based function: we mustn't clobber that different definition.

```

5565 \cs_new_protected:Npn \__str_convert:nnn #1#2#3
5566   {
5567     \cs_if_exist:cF { __str_convert_#2_#3: }
5568     {
5569       \exp_args:Nx \__str_convert:nnnn
5570         { \__str_convert_lowercase_alphanum:n {#3} }
5571         {#1} {#2} {#3}
5572     }
5573   }
5574 \cs_new_protected:Npn \__str_convert:nnnn #1#2#3#4
5575   {
5576     \cs_if_exist:cF { __str_convert_#3_#1: }

```

```

5577 {
5578   \prop_get:NnNF \g__str_alias_prop {#1} \l__str_internal_tl
5579   { \tl_set:Nn \l__str_internal_tl {#1} }
5580   \cs_if_exist:cF { __str_convert_#3_ \l__str_internal_tl : }
5581   {
5582     \file_if_exist:nTF { l3str-#2- \l__str_internal_tl .def }
5583     {
5584       \group_begin:
5585       \__str_load_catcodes:
5586       \file_input:n { l3str-#2- \l__str_internal_tl .def }
5587       \group_end:
5588     }
5589     {
5590       \tl_clear:N \l__str_internal_tl
5591       \__kernel_msg_error:nxxx { str } { unknown-#2 } {#4} {#1}
5592     }
5593   }
5594   \cs_if_exist:cF { __str_convert_#3_#1: }
5595   {
5596     \cs_gset_eq:cc { __str_convert_#3_#1: }
5597     { __str_convert_#3_ \l__str_internal_tl : }
5598   }
5599 }
5600 \cs_gset_eq:cc { __str_convert_#3_#4: } { __str_convert_#3_#1: }
5601 }

```

(End definition for __str_convert:nnn and __str_convert:nnnn.)

__str_convert_lowercase_alphanum:n
 __str_convert_lowercase_alphanum_loop:N

This function keeps only letters and digits, with upper case letters converted to lower case.

```

5602 \cs_new:Npn \__str_convert_lowercase_alphanum:n #1
5603 {
5604   \exp_after:wN \__str_convert_lowercase_alphanum_loop:N
5605   \tl_to_str:n {#1} { ? \prg_break: }
5606   \prg_break_point:
5607 }
5608 \cs_new:Npn \__str_convert_lowercase_alphanum_loop:N #1
5609 {
5610   \use_none:n #1
5611   \if_int_compare:w '#1 > 'Z \exp_stop_f:
5612   \if_int_compare:w '#1 > 'z \exp_stop_f: \else:
5613     \if_int_compare:w '#1 < 'a \exp_stop_f: \else:
5614       #1
5615     \fi:
5616   \fi:
5617   \else:
5618     \if_int_compare:w '#1 < 'A \exp_stop_f:
5619     \if_int_compare:w 1 < 1#1 \exp_stop_f:
5620       #1
5621     \fi:
5622   \else:
5623     \__str_output_byte:n { '#1 + 'a - 'A }
5624   \fi:
5625   \fi:

```

```

5626   \__str_convert_lowercase_alphanum_loop:N
5627 }

```

(End definition for `__str_convert_lowercase_alphanum:n` and `__str_convert_lowercase_alphanum_loop:N`.)

`__str_load_catcodes:` Since encoding files may be loaded at arbitrary places in a T_EX document, including within verbatim mode, we set the catcodes of all characters appearing in any encoding definition file.

```

5628 \cs_new_protected:Npn \__str_load_catcodes:
5629 {
5630   \char_set_catcode_escape:N \\\
5631   \char_set_catcode_group_begin:N \{
5632   \char_set_catcode_group_end:N \}
5633   \char_set_catcode_math_toggle:N \$
5634   \char_set_catcode_alignment:N &
5635   \char_set_catcode_parameter:N #
5636   \char_set_catcode_math_superscript:N ^
5637   \char_set_catcode_ignore:N %
5638   \char_set_catcode_space:N ~
5639   \tl_map_function:nN { abcdefghijklmnopqrstuvwxyz:ABCDEFILNPSTUX }
5640   \char_set_catcode_letter:N
5641   \tl_map_function:nN { 0123456789"?'*+-.(),'!/<>[];= }
5642   \char_set_catcode_other:N
5643   \char_set_catcode_comment:N \%
5644   \int_set:Nn \tex_endlinechar:D {32}
5645 }

```

(End definition for `__str_load_catcodes:.`)

10.3.5 Byte unescape and escape

Strings of bytes may need to be stored in auxiliary files in safe “escaping” formats. Each such escaping is only loaded as needed. By default, on input any non-byte is filtered out, while the output simply consists in letting bytes through.

`__str_filter_bytes:n` In the case of 8-bit engines, every character is a byte. For Unicode-aware engines, test the character code; non-bytes cause us to raise the flag `str_byte`. Spaces have already been given the correct category code when this function is called.

```

5646 \bool_lazy_any:nTF
5647 {
5648   \sys_if_engine luatex_p:
5649   \sys_if_engine xetex_p:
5650 }
5651 {
5652   \cs_new:Npn \__str_filter_bytes:n #1
5653   {
5654     \__str_filter_bytes_aux:N #1
5655     { ? \prg_break: }
5656     \prg_break_point:
5657   }
5658   \cs_new:Npn \__str_filter_bytes_aux:N #1
5659   {
5660     \use_none:n #1

```

```

5661         \if_int_compare:w '#1 < 256 \exp_stop_f:
5662         #1
5663         \else:
5664             \flag_raise:n { str_byte }
5665         \fi:
5666         \__str_filter_bytes_aux:N
5667     }
5668 }
5669 { \cs_new_eq:NN \__str_filter_bytes:n \use:n }

```

(End definition for `__str_filter_bytes:n` and `__str_filter_bytes_aux:N`.)

`__str_convert_unescape_:` The simplest unescaping method removes non-bytes from `\g__str_result_tl`.

```

\__str_convert_unescape_bytes:
5670 \bool_lazy_any:nTF
5671 {
5672     \sys_if_engine luatex_p:
5673     \sys_if_engine xetex_p:
5674 }
5675 {
5676     \cs_new_protected:Npn \__str_convert_unescape_:
5677     {
5678         \flag_clear:n { str_byte }
5679         \__kernel_tl_gset:Nx \g__str_result_tl
5680         { \exp_args:No \__str_filter_bytes:n \g__str_result_tl }
5681         \__str_if_flag_error:nmx { str_byte } { non-byte } { bytes }
5682     }
5683 }
5684 { \cs_new_protected:Npn \__str_convert_unescape_: { } }
5685 \cs_new_eq:NN \__str_convert_unescape_bytes: \__str_convert_unescape_:

```

(End definition for `__str_convert_unescape_:` and `__str_convert_unescape_bytes:.`)

`__str_convert_escape_:` The simplest form of escape leaves the bytes from the previous step of the conversion unchanged.

```

5686 \cs_new_protected:Npn \__str_convert_escape_: { }
5687 \cs_new_eq:NN \__str_convert_escape_bytes: \__str_convert_escape_:

```

(End definition for `__str_convert_escape_:` and `__str_convert_escape_bytes:.`)

10.3.6 Native strings

`__str_convert_decode_:` Convert each character to its character code, one at a time.

```

\__str_decode_native_char:N
5688 \cs_new_protected:Npn \__str_convert_decode_:
5689 { \__str_convert_gmap:N \__str_decode_native_char:N }
5690 \cs_new:Npn \__str_decode_native_char:N #1
5691 { #1 \s__str \int_value:w '#1 \s__str }

```

(End definition for `__str_convert_decode_:` and `__str_decode_native_char:N`.)

`__str_convert_encode_:` The conversion from an internal string to native character tokens basically maps `\char_generate:nn` through the code-points, but in non-Unicode-aware engines we use a fallback character ? rather than nothing when given a character code outside [0,255]. We detect the presence of bad characters using a flag and only produce a single error after the x-expanding assignment.

```

5692 \bool_lazy_any:nTF
5693 {
5694   \sys_if_engine luatex_p:
5695   \sys_if_engine xetex_p:
5696 }
5697 {
5698   \cs_new_protected:Npn \__str_convert_encode_:
5699     { \__str_convert_gmap_internal:N \__str_encode_native_char:n }
5700   \cs_new:Npn \__str_encode_native_char:n #1
5701     { \char_generate:nn {#1} {12} }
5702 }
5703 {
5704   \cs_new_protected:Npn \__str_convert_encode_:
5705     {
5706       \flag_clear:n { str_error }
5707       \__str_convert_gmap_internal:N \__str_encode_native_char:n
5708       \__str_if_flag_error:nx { str_error }
5709       { native-overflow } { }
5710     }
5711   \cs_new:Npn \__str_encode_native_char:n #1
5712     {
5713       \if_int_compare:w #1 > \c__str_max_byte_int
5714         \flag_raise:n { str_error }
5715       ?
5716       \else:
5717         \char_generate:nn {#1} {12}
5718       \fi:
5719     }
5720   \__kernel_msg_new:nnnn { str } { native-overflow }
5721   { Character-code-too-large-for-this-engine. }
5722   {
5723     This-engine-only-support-8-bit-characters:~
5724     valid-character-codes-are-in-the-range~[0,255].~
5725     To-manipulate-arbitrary-Unicode,-use~LuaTeX-or~XeTeX.
5726   }
5727 }

```

(End definition for __str_convert_encode_: and __str_encode_native_char:n.)

10.3.7 clist

__str_convert_decode_clist: Convert each integer to the internal form. We first turn \g__str_result_tl into a clist variable, as this avoids problems with leading or trailing commas.

```

5728 \cs_new_protected:Npn \__str_convert_decode_clist:
5729 {
5730   \clist_gset:No \g__str_result_tl \g__str_result_tl
5731   \__kernel_tl_gset:Nx \g__str_result_tl
5732   {
5733     \exp_args:No \clist_map_function:nN
5734     \g__str_result_tl \__str_decode_clist_char:n
5735   }
5736 }
5737 \cs_new:Npn \__str_decode_clist_char:n #1
5738 { #1 \s__str \int_eval:n {#1} \s__str }

```

(End definition for `__str_convert_decode_clist:` and `__str_decode_clist_char:n.`)

`__str_convert_encode_clist:` Convert the internal list of character codes to a comma-list of character codes. The first line produces a comma-list with a leading comma, removed in the next step (this also works in the empty case, since `\tl_tail:N` does not trigger an error in this case).
`__str_encode_clist_char:n`

```

5739 \cs_new_protected:Npn __str_convert_encode_clist:
5740 {
5741   __str_convert_gmap_internal:N __str_encode_clist_char:n
5742   __kernel_tl_gset:Nx \g__str_result_tl { \tl_tail:N \g__str_result_tl }
5743 }
5744 \cs_new:Npn __str_encode_clist_char:n #1 { , #1 }

```

(End definition for `__str_convert_encode_clist:` and `__str_encode_clist_char:n.`)

10.3.8 8-bit encodings

It is not clear in what situations 8-bit encodings are used, hence it is not clear what should be optimized. The current approach is reasonably efficient to convert long strings, and it scales well when using many different encodings.

The data needed to support a given 8-bit encoding is stored in a file that consists of a single function call

```

__str_declare_eight_bit_encoding:nnnn {<name>} {<modulo>} {<mapping>}
{<missing>}

```

This declares the encoding `<name>` to map bytes to Unicode characters according to the `<mapping>`, and map those bytes which are not mentioned in the `<mapping>` either to the replacement character (if they appear in `<missing>`), or to themselves. The `<mapping>` argument is a token list of pairs `{<byte>} {<Unicode>}` expressed in uppercase hexadecimal notation. The `<missing>` argument is a token list of `{<byte>}`. Every `<byte>` which does not appear in the `<mapping>` nor the `<missing>` lists maps to itself in Unicode, so for instance the `latin1` encoding has empty `<mapping>` and `<missing>` lists. The `<modulo>` is a (decimal) integer between 256 and 558 inclusive, modulo which all Unicode code points supported by the encodings must be different.

We use two integer arrays per encoding. When decoding we only use the `decode` integer array, with entry $n + 1$ (offset needed because integer array indices start at 1) equal to the Unicode code point that corresponds to the n -th byte in the encoding under consideration, or -1 if the given byte is invalid in this encoding. When encoding we use both arrays: upon seeing a code point n , we look up the entry $(1 \text{ plus } n \text{ modulo some number } M)$ in the `encode` array, which tells us the byte that might encode the given Unicode code point, then we check in the `decode` array that indeed this byte encodes the Unicode code point we want. Here, M is an encoding-dependent integer between 256 and 558 (it turns out), chosen so that among the Unicode code points that can be validly represented in the given encoding, no pair of code points have the same value modulo M .

Loop through both lists of bytes to fill in the `decode` integer array, then fill the `encode` array accordingly. For bytes that are invalid in the given encoding, store -1 in the `decode` array.

```

__str_declare_eight_bit_encoding:nnnn
__str_declare_eight_bit_aux:NNnnn
__str_declare_eight_bit_loop:Nnn
__str_declare_eight_bit_loop:Nn
5745 \cs_new_protected:Npn __str_declare_eight_bit_encoding:nnnn #1
5746 {
5747   \tl_set:Nn \l__str_internal_tl {#1}
5748   \cs_new_protected:cpn { __str_convert_decode_#1: }

```

```

5749     { \_str_convert_decode_eight_bit:n {#1} }
5750 \cs_new_protected:cpn { \_str_convert_encode_#1: }
5751     { \_str_convert_encode_eight_bit:n {#1} }
5752 \exp_args:Ncc \_str_declare_eight_bit_aux:NNnnn
5753     { g\_str_decode_#1_intarray } { g\_str_encode_#1_intarray }
5754 }
5755 \cs_new_protected:Npn \_str_declare_eight_bit_aux:NNnnn #1#2#3#4#5
5756 {
5757     \intarray_new:Nn #1 { 256 }
5758     \int_step_inline:nnn { 0 } { 255 }
5759         { \intarray_gset:Nnn #1 { 1 + ##1 } {##1} }
5760     \_str_declare_eight_bit_loop:Nnn #1
5761         #4 { \s\_str_stop \prg_break: } { }
5762     \prg_break_point:
5763     \_str_declare_eight_bit_loop:Nn #1
5764         #5 { \s\_str_stop \prg_break: }
5765     \prg_break_point:
5766     \intarray_new:Nn #2 {#3}
5767     \int_step_inline:nnn { 0 } { 255 }
5768     {
5769         \int_compare:nNnF { \intarray_item:Nn #1 { 1 + ##1 } } = { -1 }
5770         {
5771             \intarray_gset:Nnn #2
5772                 {
5773                     1 +
5774                     \int_mod:nn { \intarray_item:Nn #1 { 1 + ##1 } }
5775                     { \intarray_count:N #2 }
5776                 }
5777             {##1}
5778         }
5779     }
5780 }
5781 \cs_new_protected:Npn \_str_declare_eight_bit_loop:Nnn #1#2#3
5782 {
5783     \_str_use_none_delimit_by_s_stop:w #2 \s\_str_stop
5784     \intarray_gset:Nnn #1 { 1 + "#2 } { "#3 }
5785     \_str_declare_eight_bit_loop:Nnn #1
5786 }
5787 \cs_new_protected:Npn \_str_declare_eight_bit_loop:Nn #1#2
5788 {
5789     \_str_use_none_delimit_by_s_stop:w #2 \s\_str_stop
5790     \intarray_gset:Nnn #1 { 1 + "#2 } { -1 }
5791     \_str_declare_eight_bit_loop:Nn #1
5792 }

```

(End definition for _str_declare_eight_bit_encoding:nnnn and others.)

```

\_str_convert_decode_eight_bit:n
\_str_decode_eight_bit_aux:n
\_str_decode_eight_bit_aux:Nn

```

The map from bytes to Unicode code points is in the `decode` array corresponding to the given encoding. Define `_str_tmp:w` and pass it successively all bytes in the string. It produces an internal representation with suitable `\s_str` inserted, and the corresponding code point is obtained by looking it up in the integer array. If the entry is `-1` then issue a replacement character and raise the flag indicating that there was an error.

```

5793 \cs_new_protected:Npn \_str_convert_decode_eight_bit:n #1
5794 {

```



```

5795     \cs_set:Npx \__str_tmp:w
5796     {
5797         \exp_not:N \__str_decode_eight_bit_aux:Nn
5798         \exp_not:c { g__str_decode_#1_intarray }
5799     }
5800     \flag_clear:n { str_error }
5801     \__str_convert_gmap:N \__str_tmp:w
5802     \__str_if_flag_error:nxx { str_error } { decode-8-bit } {#1}
5803 }
5804 \cs_new:Npn \__str_decode_eight_bit_aux:Nn #1#2
5805 {
5806     #2 \s__str
5807     \exp_args:Nf \__str_decode_eight_bit_aux:n
5808     { \intarray_item:Nn #1 { 1 + '#2 } }
5809     \s__str
5810 }
5811 \cs_new:Npn \__str_decode_eight_bit_aux:n #1
5812 {
5813     \if_int_compare:w #1 < \c_zero_int
5814         \flag_raise:n { str_error }
5815         \int_value:w \c__str_replacement_char_int
5816     \else:
5817         #1
5818     \fi:
5819 }

```

(End definition for `__str_convert_decode_eight_bit:n`, `__str_decode_eight_bit_aux:n`, and `__str_decode_eight_bit_aux:Nn`.)

```

\__str_convert_encode_eight_bit:n
\__str_encode_eight_bit_aux:nnN
\__str_encode_eight_bit_aux:NNn

```

It is not practical to make an integer array with indices in the full Unicode range, so we work modulo some number, which is simply the size of the `encode` integer array for the given encoding. This gives us a candidate byte for representing a given Unicode code point. Of course taking the modulo leads to collisions so we check in the `decode` array that the byte we got is indeed correct. Otherwise the Unicode code point we started from is simply not representable in the given encoding.

```

5820 \int_new:N \l__str_modulo_int
5821 \cs_new_protected:Npn \__str_convert_encode_eight_bit:n #1
5822 {
5823     \cs_set:Npx \__str_tmp:w
5824     {
5825         \exp_not:N \__str_encode_eight_bit_aux:NNn
5826         \exp_not:c { g__str_encode_#1_intarray }
5827         \exp_not:c { g__str_decode_#1_intarray }
5828     }
5829     \flag_clear:n { str_error }
5830     \__str_convert_gmap_internal:N \__str_tmp:w
5831     \__str_if_flag_error:nxx { str_error } { encode-8-bit } {#1}
5832 }
5833 \cs_new:Npn \__str_encode_eight_bit_aux:NNn #1#2#3
5834 {
5835     \exp_args:Nf \__str_encode_eight_bit_aux:nnN
5836     {
5837         \intarray_item:Nn #1
5838         { 1 + \int_mod:nn {#3} { \intarray_count:N #1 } }

```

```

5839     }
5840     {#3}
5841     #2
5842 }
5843 \cs_new:Npn \__str_encode_eight_bit_aux:nnN #1#2#3
5844 {
5845     \int_compare:nNnTF { \intarray_item:Nn #3 { 1 + #1 } } = {#2}
5846     { \__str_output_byte:n {#1} }
5847     { \flag_raise:n { str_error } }
5848 }

```

(End definition for `__str_convert_encode_eight_bit:n`, `__str_encode_eight_bit_aux:nnN`, and `__str_encode_eight_bit_aux:NNn`.)

10.4 Messages

General messages, and messages for the encodings and escapings loaded by default (“native”, and “bytes”).

```

5849 \__kernel_msg_new:nnn { str } { unknown-esc }
5850 { Escaping-scheme~'#1'~(filtered:~'#2')~unknown. }
5851 \__kernel_msg_new:nnn { str } { unknown-enc }
5852 { Encoding-scheme~'#1'~(filtered:~'#2')~unknown. }
5853 \__kernel_msg_new:nnnn { str } { native-escaping }
5854 { The~'native'~encoding-scheme~does~not~support~any~escaping. }
5855 {
5856     Since~native~strings~do~not~consist~in~bytes,~
5857     none~of~the~escaping~methods~make~sense.~
5858     The~specified~escaping,~'#1',~will be ignored.
5859 }
5860 \__kernel_msg_new:nnn { str } { file-not-found }
5861 { File~'l3str-#1.def'~not~found. }

```

Message used when the “bytes” unescaping fails because the string given to `\str_set_convert:Nnnn` contains a non-byte. This cannot happen for the -8-bit engines. Messages used for other escapings and encodings are defined in each definition file.

```

5862 \bool_lazy_any:nT
5863 {
5864     \sys_if_engine luatex_p:
5865     \sys_if_engine xetex_p:
5866 }
5867 {
5868     \__kernel_msg_new:nnnn { str } { non-byte }
5869     { String~invalid~in~escaping~'#1':~it~may~only~contain~bytes. }
5870     {
5871         Some~characters~in~the~string~you~asked~to~convert~are~not~
5872         8-bit~characters.~Perhaps~the~string~is~a~'native'~Unicode~string?~
5873         If~it~is,~try~using\\
5874         \\
5875         \iow_indent:n
5876         {
5877             \iow_char:N\\str_set_convert:Nnnn \\
5878             \ \ <str-var>~\{<string>~\}~\{~native~\}~\{<target-encoding>~\}
5879         }
5880     }
5881 }

```

Those messages are used when converting to and from 8-bit encodings.

```

5882 \__kernel_msg_new:nnnn { str } { decode-8-bit }
5883 { Invalid-string-in-encoding-’#1’. }
5884 {
5885     LaTeX-came-across-a-byte-which-is-not-defined-to-represent~
5886     any-character-in-the-encoding-’#1’.
5887 }
5888 \__kernel_msg_new:nnnn { str } { encode-8-bit }
5889 { Unicode-string-cannot-be-converted-to-encoding-’#1’. }
5890 {
5891     The-encoding-’#1’-only-contains-a-subset-of-all-Unicode-characters.~
5892     LaTeX-was-asked-to-convert-a-string-to-that-encoding,~but~that~
5893     string-contains-a-character-that-’#1’~does-not-support.
5894 }

```

10.5 Escaping definitions

Several of those encodings are defined by the pdf file format. The following byte storage methods are defined:

- **bytes** (default), non-bytes are filtered out, and bytes are left untouched (this is defined by default);
- **hex** or **hexadecimal**, as per the pdfTeX primitive `\pdfescapehex`
- **name**, as per the pdfTeX primitive `\pdfescapename`
- **string**, as per the pdfTeX primitive `\pdfescapestring`
- **url**, as per the percent encoding of urls.

10.5.1 Unescape methods

`__str_convert_unescape_hex:` Take chars two by two, and interpret each pair as the hexadecimal code for a byte. Anything else than hexadecimal digits is ignored, raising the flag. A string which contains an odd number of hexadecimal digits gets 0 appended to it: this is equivalent to appending a 0 in all cases, and dropping it if it is alone.

```

5895 \cs_new_protected:Npn \__str_convert_unescape_hex:
5896 {
5897     \group_begin:
5898     \flag_clear:n { str_error }
5899     \int_set:Nn \tex_escapechar:D { 92 }
5900     \__kernel_tl_gset:Nx \g__str_result_tl
5901     {
5902         \__str_output_byte:w "
5903         \exp_last_unbraced:Nf \__str_unescape_hex_auxi:N
5904         { \tl_to_str:N \g__str_result_tl }
5905         0 { ? 0 - 1 \prg_break: }
5906         \prg_break_point:
5907         \__str_output_end:
5908     }
5909     \__str_if_flag_error:nnx { str_error } { unescape-hex } { }
5910     \group_end:
5911 }

```

```

5912 \cs_new:Npn \__str_unescape_hex_auxi:N #1
5913 {
5914   \use_none:n #1
5915   \__str_hexadecimal_use:NTF #1
5916   { \__str_unescape_hex_auxii:N }
5917   {
5918     \flag_raise:n { str_error }
5919     \__str_unescape_hex_auxi:N
5920   }
5921 }
5922 \cs_new:Npn \__str_unescape_hex_auxii:N #1
5923 {
5924   \use_none:n #1
5925   \__str_hexadecimal_use:NTF #1
5926   {
5927     \__str_output_end:
5928     \__str_output_byte:w " \__str_unescape_hex_auxi:N
5929   }
5930   {
5931     \flag_raise:n { str_error }
5932     \__str_unescape_hex_auxii:N
5933   }
5934 }
5935 \__kernel_msg_new:nnnn { str } { unescape-hex }
5936 { String~invalid~in~escaping~'hex':~only~hexadecimal~digits~allowed. }
5937 {
5938   Some~characters~in~the~string~you~asked~to~convert~are~not~
5939   hexadecimal~digits~(0-9,~A-F,~a-f)~nor~spaces.
5940 }

```

(End definition for __str_convert_unescape_hex:, __str_unescape_hex_auxi:N, and __str_unescape_hex_auxii:N.)

<pre> __str_convert_unescape_name: __str_unescape_name_loop:wNN __str_convert_unescape_url: __str_unescape_url_loop:wNN </pre>	<p>The __str_convert_unescape_name: function replaces each occurrence of # followed by two hexadecimal digits in \g__str_result_tl by the corresponding byte. The url function is identical, with escape character % instead of #. Thus we define the two together. The arguments of __str_tmp:w are the character code of # or % in hexadecimal, the name of the main function to define, and the name of the auxiliary which performs the loop.</p>
--	---

The looping auxiliary #3 finds the next escape character, reads the following two characters, and tests them. The test __str_hexadecimal_use:NTF leaves the upper-case digit in the input stream, hence we surround the test with __str_output_byte:w " and __str_output_end:. If both characters are hexadecimal digits, they should be removed before looping: this is done by \use_i:nnn. If one of the characters is not a hexadecimal digit, then feed "#1 to __str_output_byte:w to produce the escape character, raise the flag, and call the looping function followed by the two characters (remove \use_i:nnn).

```

5941 \cs_set_protected:Npn \__str_tmp:w #1#2#3
5942 {
5943   \cs_new_protected:cpn { __str_convert_unescape_#2: }
5944   {
5945     \group_begin:
5946     \flag_clear:n { str_byte }

```

```

5947 \flag_clear:n { str_error }
5948 \int_set:Nn \tex_escapechar:D { 92 }
5949 \__kernel_tl_gset:Nx \g__str_result_tl
5950 {
5951 \exp_after:wN #3 \g__str_result_tl
5952 #1 ? { ? \prg_break: }
5953 \prg_break_point:
5954 }
5955 \__str_if_flag_error:nxx { str_byte } { non-byte } { #2 }
5956 \__str_if_flag_error:nxx { str_error } { unescape-#2 } { }
5957 \group_end:
5958 }
5959 \cs_new:Npn #3 ##1#1##2##3
5960 {
5961 \__str_filter_bytes:n {##1}
5962 \use_none:n ##3
5963 \__str_output_byte:w "
5964 \__str_hexadecimal_use:NTF ##2
5965 {
5966 \__str_hexadecimal_use:NTF ##3
5967 { }
5968 {
5969 \flag_raise:n { str_error }
5970 * 0 + '#1 \use_i:nn
5971 }
5972 }
5973 {
5974 \flag_raise:n { str_error }
5975 0 + '#1 \use_i:nn
5976 }
5977 \__str_output_end:
5978 \use_i:nnn #3 ##2##3
5979 }
5980 \__kernel_msg_new:nnnn { str } { unescape-#2 }
5981 { String~invalid~in~escaping~'#2'. }
5982 {
5983 LaTeX~came~across~the~escape~character~'#1'~not~followed~by~
5984 two~hexadecimal~digits.~This~is~invalid~in~the~escaping~'#2'.
5985 }
5986 }
5987 \exp_after:wN \__str_tmp:w \c_hash_str { name }
5988 \__str_unescape_name_loop:wNN
5989 \exp_after:wN \__str_tmp:w \c_percent_str { url }
5990 \__str_unescape_url_loop:wNN

```

(End definition for `__str_convert_unescape_name:` and others.)

```

\__str_convert_unescape_string:
\__str_unescape_string_newlines:wN
\__str_unescape_string_loop:wNNN
\__str_unescape_string_repeat:NNNNNN

```

The **string** escaping is somewhat similar to the **name** and **url** escapings, with escape character `\`. The first step is to convert all three line endings, `^^J`, `^^M`, and `^^M^^J` to the common `^^J`, as per the PDF specification. This step cannot raise the flag.

Then the following escape sequences are decoded.

`\n` Line feed (10)

`\r` Carriage return (13)

`\t` Horizontal tab (9)

`\b` Backspace (8)

`\f` Form feed (12)

`\(` Left parenthesis

`\)` Right parenthesis

`\\` Backslash

`\ddd` (backslash followed by 1 to 3 octal digits) Byte `ddd` (octal), subtracting 256 in case of overflow.

If followed by an end-of-line character, the backslash and the end-of-line are ignored. If followed by anything else, the backslash is ignored, raising the error flag.

```
5991 \group_begin:
5992   \char_set_catcode_other:N \^^J
5993   \char_set_catcode_other:N \^^M
5994   \cs_set_protected:Npn \__str_tmp:w #1
5995   {
5996     \cs_new_protected:Npn \__str_convert_unescape_string:
5997     {
5998       \group_begin:
5999         \flag_clear:n { str_byte }
6000         \flag_clear:n { str_error }
6001         \int_set:Nn \tex_escapechar:D { 92 }
6002         \__kernel_tl_gset:Nx \g__str_result_tl
6003         {
6004           \exp_after:wN \__str_unescape_string_newlines:wN
6005           \g__str_result_tl \prg_break: ^^M ?
6006           \prg_break_point:
6007         }
6008         \__kernel_tl_gset:Nx \g__str_result_tl
6009         {
6010           \exp_after:wN \__str_unescape_string_loop:wNNN
6011           \g__str_result_tl #1 ?? { ? \prg_break: }
6012           \prg_break_point:
6013         }
6014         \__str_if_flag_error:nnx { str_byte } { non-byte } { string }
6015         \__str_if_flag_error:nnx { str_error } { unescape-string } { }
6016       \group_end:
6017     }
6018   }
6019   \exp_args:No \__str_tmp:w { \c_backslash_str }
6020   \exp_last_unbraced:NNNNo
6021   \cs_new:Npn \__str_unescape_string_loop:wNNN #1 \c_backslash_str #2#3#4
6022   {
6023     \__str_filter_bytes:n {#1}
6024     \use_none:n #4
6025     \__str_output_byte:w '
6026     \__str_octal_use:NTF #2
6027     {
6028       \__str_octal_use:NTF #3
```

```

6029         {
6030             \__str_octal_use:NTF #4
6031             {
6032                 \if_int_compare:w #2 > 3 \exp_stop_f:
6033                 - 256
6034                 \fi:
6035                 \__str_unescape_string_repeat:NNNNNN
6036             }
6037             { \__str_unescape_string_repeat:NNNNNN ? }
6038         }
6039         { \__str_unescape_string_repeat:NNNNNN ?? }
6040     }
6041     {
6042         \str_case_e:nnF {#2}
6043         {
6044             { \c_backslash_str } { 134 }
6045             { ( } { 50 }
6046             { ) } { 51 }
6047             { r } { 15 }
6048             { f } { 14 }
6049             { n } { 12 }
6050             { t } { 11 }
6051             { b } { 10 }
6052             { ^^J } { 0 - 1 }
6053         }
6054         {
6055             \flag_raise:n { str_error }
6056             0 - 1 \use_i:nn
6057         }
6058     }
6059     \__str_output_end:
6060     \use_i:nn \__str_unescape_string_loop:wNNN #2#3#4
6061 }
6062 \cs_new:Npn \__str_unescape_string_repeat:NNNNNN #1#2#3#4#5#6
6063 { \__str_output_end: \__str_unescape_string_loop:wNNN }
6064 \cs_new:Npn \__str_unescape_string_newlines:wN #1 ^^M #2
6065 {
6066     #1
6067     \if_charcode:w ^^J #2 \else: ^^J \fi:
6068     \__str_unescape_string_newlines:wN #2
6069 }
6070 \__kernel_msg_new:nnnn { str } { unescape-string }
6071 { String~invalid~in~escaping~'string'. }
6072 {
6073     LaTeX~came~across~an~escape~character~'\c_backslash_str'~
6074     not~followed~by~any~of:~'n',~'r',~'t',~'b',~'f',~'(',~')',~
6075     '\c_backslash_str',~one~to~three~octal~digits,~or~the~end~
6076     of~a~line.
6077 }
6078 \group_end:

```

(End definition for __str_convert_unescape_string: and others.)

10.5.2 Escape methods

Currently, none of the escape methods can lead to errors, assuming that their input is made out of bytes.

`__str_convert_escape_hex:` Loop and convert each byte to hexadecimal.

```
\__str_escape_hex_char:N
6079 \cs_new_protected:Npn \__str_convert_escape_hex:
6080 { \__str_convert_gmap:N \__str_escape_hex_char:N }
6081 \cs_new:Npn \__str_escape_hex_char:N #1
6082 { \__str_output_hexadecimal:n { '#1 } }
```

(End definition for `__str_convert_escape_hex:` and `__str_escape_hex_char:N`.)

`__str_convert_escape_name:` For each byte, test whether it should be output as is, or be “hash-encoded”. Roughly, bytes outside the range [“2A”, “7E”] are hash-encoded. We keep two lists of exceptions: `__str_escape_name_char:n` characters in `\c__str_escape_name_not_str` are not hash-encoded, and characters in the `\c__str_escape_name_str` are encoded.

```
\__str_escape_name_char:n
\__str_if_escape_name:nTF
\c__str_escape_name_str
\c__str_escape_name_not_str
6083 \str_const:Nn \c__str_escape_name_not_str { ! " $ & ' } %$
6084 \str_const:Nn \c__str_escape_name_str { { } / < > [ ] }
6085 \cs_new_protected:Npn \__str_convert_escape_name:
6086 { \__str_convert_gmap:N \__str_escape_name_char:n }
6087 \cs_new:Npn \__str_escape_name_char:n #1
6088 {
6089   \__str_if_escape_name:nTF {#1} {#1}
6090   { \c_hash_str \__str_output_hexadecimal:n { '#1 } }
6091 }
6092 \prg_new_conditional:Npnn \__str_if_escape_name:n #1 { TF }
6093 {
6094   \if_int_compare:w '#1 < "2A \exp_stop_f:
6095   \__str_if_contains_char:NnTF \c__str_escape_name_not_str {#1}
6096   \prg_return_true: \prg_return_false:
6097   \else:
6098   \if_int_compare:w '#1 > "7E \exp_stop_f:
6099   \prg_return_false:
6100   \else:
6101   \__str_if_contains_char:NnTF \c__str_escape_name_str {#1}
6102   \prg_return_false: \prg_return_true:
6103   \fi:
6104   \fi:
6105 }
```

(End definition for `__str_convert_escape_name:` and others.)

`__str_convert_escape_string:` Any character below (and including) space, and any character above (and including) `del`, are converted to octal. One backslash is added before each parenthesis and backslash.

```
\__str_escape_string_char:N
\__str_if_escape_string:NnTF
\c__str_escape_string_str
6106 \str_const:Nx \c__str_escape_string_str
6107 { \c_backslash_str ( ) }
6108 \cs_new_protected:Npn \__str_convert_escape_string:
6109 { \__str_convert_gmap:N \__str_escape_string_char:N }
6110 \cs_new:Npn \__str_escape_string_char:N #1
6111 {
6112   \__str_if_escape_string:NnTF #1
6113   {
6114     \__str_if_contains_char:NnT
```



```

6115         \c__str_escape_string_str {#1}
6116         { \c_backslash_str }
6117     #1
6118 }
6119 {
6120     \c_backslash_str
6121     \int_div_truncate:nn {'#1} {64}
6122     \int_mod:nn { \int_div_truncate:nn {'#1} { 8 } } { 8 }
6123     \int_mod:nn {'#1} { 8 }
6124 }
6125 }
6126 \prg_new_conditional:Npnn \__str_if_escape_string:N #1 { TF }
6127 {
6128     \if_int_compare:w '#1 < "21 \exp_stop_f:
6129     \prg_return_false:
6130 }else:
6131     \if_int_compare:w '#1 > "7E \exp_stop_f:
6132     \prg_return_false:
6133 }else:
6134     \prg_return_true:
6135 }fi:
6136 }fi:
6137 }

```

(End definition for __str_convert_escape_string: and others.)

__str_convert_escape_url: This function is similar to __str_convert_escape_name:, escaping different characters.

```

\__str_escape_url_char:n
\__str_if_escape_url:nTF
6138 \cs_new_protected:Npn \__str_convert_escape_url:
6139 { \__str_convert_gmap:N \__str_escape_url_char:n }
6140 \cs_new:Npn \__str_escape_url_char:n #1
6141 {
6142     \__str_if_escape_url:nTF {#1} {#1}
6143     { \c_percent_str \__str_output_hexadecimal:n { '#1 } }
6144 }
6145 \prg_new_conditional:Npnn \__str_if_escape_url:n #1 { TF }
6146 {
6147     \if_int_compare:w '#1 < "41 \exp_stop_f:
6148     \__str_if_contains_char:nnTF { "-.<> } {#1}
6149     \prg_return_true: \prg_return_false:
6150 }else:
6151     \if_int_compare:w '#1 > "7E \exp_stop_f:
6152     \prg_return_false:
6153 }else:
6154     \__str_if_contains_char:nnTF { [ ] } {#1}
6155     \prg_return_false: \prg_return_true:
6156 }fi:
6157 }fi:
6158 }

```

(End definition for __str_convert_escape_url:, __str_escape_url_char:n, and __str_if_escape_url:nTF.)

10.6 Encoding definitions

The **native** encoding is automatically defined. Other encodings are loaded as needed. The following encodings are supported:

- UTF-8;
- UTF-16, big-, little-endian, or with byte order mark;
- UTF-32, big-, little-endian, or with byte order mark;
- the ISO 8859 code pages, numbered from 1 to 16, skipping the inexistent ISO 8859-12.

10.6.1 utf-8 support

```
__str_convert_encode_utf8: Loop through the internal string, and convert each character to its UTF-8 representation.
    __str_encode_utf_viii_char:n The representation is built from the right-most (least significant) byte to the left-most
    __str_encode_utf_viii_loop:wnnw (most significant) byte. Continuation bytes are in the range [128, 191], taking 64 different
values, hence we roughly want to express the character code in base 64, shifting the
first digit in the representation by some number depending on how many continuation
bytes there are. In the range [0, 127], output the corresponding byte directly. In the
range [128, 2047], output the remainder modulo 64, plus 128 as a continuation byte, then
output the quotient (which is in the range [0, 31]), shifted by 192. In the next range,
[2048, 65535], split the character code into residue and quotient modulo 64, output the
residue as a first continuation byte, then repeat; this leaves us with a quotient in the
range [0, 15], which we output shifted by 224. The last range, [65536, 1114111], follows
the same pattern: once we realize that dividing twice by 64 leaves us with a number
larger than 15, we repeat, producing a last continuation byte, and offset the quotient by
240 for the leading byte.
```

How is that implemented? `__str_encode_utf_viii_loop:wnnw` takes successive quotients as its first argument, the quotient from the previous step as its second argument (except in step 1), the bound for quotients that trigger one more step or not, and finally the offset used if this step should produce the leading byte. Leading bytes can be in the ranges [0, 127], [192, 223], [224, 239], and [240, 247] (really, that last limit should be 244 because Unicode stops at the code point 1114111). At each step, if the quotient `#1` is less than the limit `#3` for that range, output the leading byte (`#1` shifted by `#4`) and stop. Otherwise, we need one more step: use the quotient of `#1` by 64, and `#1` as arguments for the looping auxiliary, and output the continuation byte corresponding to the remainder `#2 - 64#1 + 128`. The bizarre construction `- 1 + 0 *` removes the spurious initial continuation byte (better methods welcome).

```
6159 \cs_new_protected:cpn { __str_convert_encode_utf8: }
6160 { __str_convert_gmap_internal:N __str_encode_utf_viii_char:n }
6161 \cs_new:Npn __str_encode_utf_viii_char:n #1
6162 {
6163     __str_encode_utf_viii_loop:wnnw #1 ; - 1 + 0 * ;
6164     { 128 } { 0 }
6165     { 32 } { 192 }
6166     { 16 } { 224 }
6167     { 8 } { 240 }
6168     \s__str_stop
6169 }
6170 \cs_new:Npn __str_encode_utf_viii_loop:wnnw #1; #2; #3#4 #5 \s__str_stop
```

```

6171 {
6172   \if_int_compare:w #1 < #3 \exp_stop_f:
6173     \__str_output_byte:n { #1 + #4 }
6174     \exp_after:wN \__str_use_none_delimit_by_s_stop:w
6175   \fi:
6176   \exp_after:wN \__str_encode_utf_viii_loop:wnnw
6177     \int_value:w \int_div_truncate:nn {#1} {64} ; #1 ;
6178     #5 \s__str_stop
6179   \__str_output_byte:n { #2 - 64 * ( #1 - 2 ) }
6180 }

```

(End definition for __str_convert_encode_utf8:, __str_encode_utf_viii_char:n, and __str_encode_utf_viii_loop:wnnw.)

\l__str_missing_flag When decoding a string that is purportedly in the UTF-8 encoding, four different errors can occur, signalled by a specific flag for each (we define those flags using \flag_clear_new:n rather than \flag_new:n, because they are shared with other encoding definition files).

- “Missing continuation byte”: a leading byte is not followed by the right number of continuation bytes.
- “Extra continuation byte”: a continuation byte appears where it was not expected, *i.e.*, not after an appropriate leading byte.
- “Overlong”: a Unicode character is expressed using more bytes than necessary, for instance, "C0"80 for the code point 0, instead of a single null byte.
- “Overflow”: this occurs when decoding produces Unicode code points greater than 1114111.

We only raise one L^AT_EX3 error message, combining all the errors which occurred. In the short message, the leading comma must be removed to get a grammatically correct sentence. In the long text, first remind the user what a correct UTF-8 string should look like, then add error-specific information.

```

6181 \flag_clear_new:n { str_missing }
6182 \flag_clear_new:n { str_extra }
6183 \flag_clear_new:n { str_overlong }
6184 \flag_clear_new:n { str_overflow }
6185 \__kernel_msg_new:nnnn { str } { utf8-decode }
6186 {
6187   Invalid~UTF-8~string:
6188   \exp_last_unbraced:Nf \use_none:n
6189   {
6190     \__str_if_flag_times:nT { str_missing } { ,~missing~continuation~byte }
6191     \__str_if_flag_times:nT { str_extra } { ,~extra~continuation~byte }
6192     \__str_if_flag_times:nT { str_overlong } { ,~overlong~form }
6193     \__str_if_flag_times:nT { str_overflow } { ,~code~point~too~large }
6194   }
6195   .
6196 }
6197 {
6198   In~the~UTF-8~encoding,~each~Unicode~character~consists~in~
6199   1~to~4~bytes,~with~the~following~bit~pattern: \\\
6200   \iow_indent:n

```

```

6201     {
6202         Code-point~\ \ \ \ <~128:~0xxxxxxx \\\
6203         Code-point~\ \ \ \ <~2048:~110xxxxx~10xxxxxx \\\
6204         Code-point~\ \ \ \ <~65536:~1110xxxx~10xxxxxx~10xxxxxx \\\
6205         Code-point~ \ \ \ \ <~1114112:~11110xxx~10xxxxxx~10xxxxxx~10xxxxxx \\\
6206     }
6207 Bytes-of-the-form-10xxxxxx-are-called-continuation-bytes.
6208 \flag_if_raised:nT { str_missing }
6209 {
6210     \\\
6211     A-leading-byte~(in-the-range~[192,255])~was~not~followed-by~
6212     the-appropriate-number-of~continuation-bytes.
6213 }
6214 \flag_if_raised:nT { str_extra }
6215 {
6216     \\\
6217     LaTeX-came-across-a-continuation-byte-when-it-was-not-expected.
6218 }
6219 \flag_if_raised:nT { str_overlong }
6220 {
6221     \\\
6222     Every~Unicode-code-point~must~be~expressed~in~the~shortest~
6223     possible~form.~For~instance,~'0xC0'~'0x83'~is~not~a~valid~
6224     representation~for~the~code~point~3.
6225 }
6226 \flag_if_raised:nT { str_overflow }
6227 {
6228     \\\
6229     Unicode-limits-code-points-to-the-range~[0,1114111].
6230 }
6231 }

```

(End definition for `\l__str_missing_flag` and others.)

`__str_convert_decode_utf8:` Decoding is significantly harder than encoding. As before, lower some flags, which are tested at the end (in bulk, to trigger at most one L^AT_EX3 error, as explained above). We expect successive multi-byte sequences of the form $\langle start\ byte \rangle \langle continuation\ bytes \rangle$. The `_start` auxiliary tests the first byte:

- [0, "7F]: the byte stands alone, and is converted to its own character code;
- ["80, "BF]: unexpected continuation byte, raise the appropriate flag, and convert that byte to the replacement character "FFFD;
- ["C0, "FF]: this byte should be followed by some continuation byte(s).

In the first two cases, `\use_none_delimit_by_q_stop:w` removes data that only the third case requires, namely the limits of ranges of Unicode characters which can be expressed with 1, 2, 3, or 4 bytes.

We can now concentrate on the multi-byte case and the `_continuation` auxiliary. We expect `#3` to be in the range ["80, "BF]. The test for this goes as follows: if the character code is less than "80, we compare it to `—"C0`, yielding `false`; otherwise to `"C0`, yielding `true` in the range ["80, "BF] and `false` otherwise. If we find that the byte is not a continuation range, stop the current slew of bytes, output the replacement

character, and continue parsing with the `_start` auxiliary, starting at the byte we just tested. Once we know that the byte is a continuation byte, leave it behind us in the input stream, compute what code point the bytes read so far would produce, and feed that number to the `_aux` function.

The `_aux` function tests whether we should look for more continuation bytes or not. If the number it receives as `#1` is less than the maximum `#4` for the current range, then we are done: check for an overlong representation by comparing `#1` with the maximum `#3` for the previous range. Otherwise, we call the `_continuation` auxiliary again, after shifting the “current code point” by `#4` (maximum from the range we just checked).

Two additional tests are needed: if we reach the end of the list of range maxima and we are still not done, then we are faced with an overflow. Clean up, and again insert the code point "FFFD" for the replacement character. Also, every time we read a byte, we need to check whether we reached the end of the string. In a correct UTF-8 string, this happens automatically when the `_start` auxiliary leaves its first argument in the input stream: the end-marker begins with `\prg_break:`, which ends the loop. On the other hand, if the end is reached when looking for a continuation byte, the `\use_none:n #3` construction removes the first token from the end-marker, and leaves the `_end` auxiliary, which raises the appropriate error flag before ending the mapping.

```

6232 \cs_new_protected:cpn { __str_convert_decode_utf8: }
6233 {
6234   \flag_clear:n { str_error }
6235   \flag_clear:n { str_missing }
6236   \flag_clear:n { str_extra }
6237   \flag_clear:n { str_overlong }
6238   \flag_clear:n { str_overflow }
6239   \__kernel_tl_gset:Nx \g__str_result_tl
6240   {
6241     \exp_after:wN \__str_decode_utf_viii_start:N \g__str_result_tl
6242     { \prg_break: \__str_decode_utf_viii_end: }
6243     \prg_break_point:
6244   }
6245   \__str_if_flag_error:nxx { str_error } { utf8-decode } { }
6246 }
6247 \cs_new:Npn \__str_decode_utf_viii_start:N #1
6248 {
6249   #1
6250   \if_int_compare:w '#1 < "C0 \exp_stop_f:
6251     \s__str
6252     \if_int_compare:w '#1 < "80 \exp_stop_f:
6253       \int_value:w '#1
6254     \else:
6255       \flag_raise:n { str_extra }
6256       \flag_raise:n { str_error }
6257       \int_use:N \c__str_replacement_char_int
6258     \fi:
6259   \else:
6260     \exp_after:wN \__str_decode_utf_viii_continuation:wwN
6261     \int_value:w \int_eval:n { '#1 - "C0 } \exp_after:wN
6262   \fi:
6263   \s__str
6264   \__str_use_none_delimit_by_s_stop:w {"80} {"800} {"10000} {"110000} \s__str_stop
6265   \__str_decode_utf_viii_start:N

```

```

6266     }
6267 \cs_new:Npn \__str_decode_utf_viii_continuation:wwN
6268     #1 \s__str #2 \__str_decode_utf_viii_start:N #3
6269     {
6270     \use_none:n #3
6271     \if_int_compare:w '#3 <
6272         \if_int_compare:w '#3 < "80 \exp_stop_f: - \fi:
6273         "C0 \exp_stop_f:
6274         #3
6275         \exp_after:wN \__str_decode_utf_viii_aux:wNnnwN
6276         \int_value:w \int_eval:n { #1 * "40 + '#3 - "80 } \exp_after:wN
6277     \else:
6278         \s__str
6279         \flag_raise:n { str_missing }
6280         \flag_raise:n { str_error }
6281         \int_use:N \c__str_replacement_char_int
6282     \fi:
6283     \s__str
6284     #2
6285     \__str_decode_utf_viii_start:N #3
6286     }
6287 \cs_new:Npn \__str_decode_utf_viii_aux:wNnnwN
6288     #1 \s__str #2#3#4 #5 \__str_decode_utf_viii_start:N #6
6289     {
6290     \if_int_compare:w #1 < #4 \exp_stop_f:
6291         \s__str
6292         \if_int_compare:w #1 < #3 \exp_stop_f:
6293             \flag_raise:n { str_overlong }
6294             \flag_raise:n { str_error }
6295             \int_use:N \c__str_replacement_char_int
6296         \else:
6297             #1
6298         \fi:
6299     \else:
6300         \if_meaning:w \s__str_stop #5
6301             \__str_decode_utf_viii_overflow:w #1
6302         \fi:
6303         \exp_after:wN \__str_decode_utf_viii_continuation:wwN
6304         \int_value:w \int_eval:n { #1 - #4 } \exp_after:wN
6305     \fi:
6306     \s__str
6307     #2 {#4} #5
6308     \__str_decode_utf_viii_start:N
6309     }
6310 \cs_new:Npn \__str_decode_utf_viii_overflow:w #1 \fi: #2 \fi:
6311     {
6312     \fi: \fi:
6313     \flag_raise:n { str_overflow }
6314     \flag_raise:n { str_error }
6315     \int_use:N \c__str_replacement_char_int
6316     }
6317 \cs_new:Npn \__str_decode_utf_viii_end:
6318     {
6319     \s__str

```

```

6320     \flag_raise:n { str_missing }
6321     \flag_raise:n { str_error }
6322     \int_use:N \c__str_replacement_char_int \s__str
6323     \prg_break:
6324 }

```

(End definition for `__str_convert_decode_utf8:` and others.)

10.6.2 utf-16 support

The definitions are done in a category code regime where the bytes 254 and 255 used by the byte order mark have catcode 12.

```

6325 \group_begin:
6326   \char_set_catcode_other:N ^^fe
6327   \char_set_catcode_other:N ^^ff

```

`__str_convert_encode_utf16:` When the endianness is not specified, it is big-endian by default, and we add a byte-order mark. Convert characters one by one in a loop, with different behaviours depending on the character code.

`__str_encode_utf_xvi_aux:N`
`__str_encode_utf_xvi_char:n`

- [0, "D7FF]: converted to two bytes;
- ["D800, "DFFF] are used as surrogates: they cannot be converted and are replaced by the replacement character;
- ["E000, "FFFF]: converted to two bytes;
- ["10000, "10FFFF]: converted to a pair of surrogates, each two bytes. The magic "D7C0 is "D800 – "10000/"400.

For the duration of this operation, `__str_tmp:w` is defined as a function to convert a number in the range [0, "FFFF] to a pair of bytes (either big endian or little endian), by feeding the quotient of the division of #1 by "100, followed by #1 to `__str_encode_utf_xvi_be:nn` or its `le` analog: those compute the remainder, and output two bytes for the quotient and remainder.

```

6328   \cs_new_protected:cpn { __str_convert_encode_utf16: }
6329   {
6330     \__str_encode_utf_xvi_aux:N \__str_output_byte_pair_be:n
6331     \tl_gput_left:Nx \g__str_result_tl { ^^fe ^^ff }
6332   }
6333   \cs_new_protected:cpn { __str_convert_encode_utf16be: }
6334   { \__str_encode_utf_xvi_aux:N \__str_output_byte_pair_be:n }
6335   \cs_new_protected:cpn { __str_convert_encode_utf16le: }
6336   { \__str_encode_utf_xvi_aux:N \__str_output_byte_pair_le:n }
6337   \cs_new_protected:Npn \__str_encode_utf_xvi_aux:N #1
6338   {
6339     \flag_clear:n { str_error }
6340     \cs_set_eq:NN \__str_tmp:w #1
6341     \__str_convert_gmap_internal:N \__str_encode_utf_xvi_char:n
6342     \__str_if_flag_error:nxx { str_error } { utf16-encode } { }
6343   }
6344   \cs_new:Npn \__str_encode_utf_xvi_char:n #1
6345   {
6346     \if_int_compare:w #1 < "D800 \exp_stop_f:

```

```

6347     \__str_tmp:w {#1}
6348 \else:
6349     \if_int_compare:w #1 < "10000 \exp_stop_f:
6350     \if_int_compare:w #1 < "E000 \exp_stop_f:
6351     \flag_raise:n { str_error }
6352     \__str_tmp:w { \c__str_replacement_char_int }
6353 \else:
6354     \__str_tmp:w {#1}
6355 \fi:
6356 \else:
6357     \exp_args:Nf \__str_tmp:w { \int_div_truncate:nn {#1} {"400} + "D7C0 }
6358     \exp_args:Nf \__str_tmp:w { \int_mod:nn {#1} {"400} + "DC00 }
6359 \fi:
6360 \fi:
6361 }

```

(End definition for __str_convert_encode_utf16: and others.)

\l__str_missing_flag When encoding a Unicode string to UTF-16, only one error can occur: code points in the range ["D800, "DFFF], corresponding to surrogates, cannot be encoded. We use the
 \l__str_extra_flag all-purpose flag @@_error to signal that error.
 \l__str_end_flag

When decoding a Unicode string which is purportedly in UTF-16, three errors can occur: a missing trail surrogate, an unexpected trail surrogate, and a string containing an odd number of bytes.

```

6362 \flag_clear_new:n { str_missing }
6363 \flag_clear_new:n { str_extra }
6364 \flag_clear_new:n { str_end }
6365 \__kernel_msg_new:nnnn { str } { utf16-encode }
6366 { Unicode-string-cannot-be-expressed-in-UTF-16:~surrogate. }
6367 {
6368     Surrogate-code-points~(in-the-range~[U+D800,~U+DFFF])~
6369     can-be-expressed-in-the~UTF-8-and-UTF-32-encodings,~
6370     but-not-in-the-UTF-16-encoding.
6371 }
6372 \__kernel_msg_new:nnnn { str } { utf16-decode }
6373 {
6374     Invalid-UTF-16-string:
6375     \exp_last_unbraced:Nf \use_none:n
6376     {
6377         \__str_if_flag_times:nT { str_missing } { ,~missing~trail~surrogate }
6378         \__str_if_flag_times:nT { str_extra } { ,~extra~trail~surrogate }
6379         \__str_if_flag_times:nT { str_end } { ,~odd~number~of~bytes }
6380     }
6381     .
6382 }
6383 {
6384     In~the~UTF-16-encoding,~each~Unicode~character~is~encoded~as~
6385     2-or-4-bytes: \\\
6386     \iow_indent:n
6387     {
6388         Code-point-in~[U+0000,~U+D7FF]:~two-bytes \\\
6389         Code-point-in~[U+D800,~U+DFFF]:~illegal \\\
6390         Code-point-in~[U+E000,~U+FFFF]:~two-bytes \\\
6391         Code-point-in~[U+10000,~U+10FFFF]:~

```



```

6392         a~lead-surrogate-and-a-trail-surrogate \\
6393     }
6394     Lead-surrogates-are-pairs-of-bytes-in-the-range-[0xD800,~0xDBFF],~
6395     and-trail-surrogates-are-in-the-range-[0xDC00,~0xDFFF].
6396     \flag_if_raised:nT { str_missing }
6397     {
6398         \\
6399         A~lead-surrogate-was-not-followed-by-a-trail-surrogate.
6400     }
6401     \flag_if_raised:nT { str_extra }
6402     {
6403         \\
6404         LaTeX-came-across-a-trail-surrogate-when-it-was-not-expected.
6405     }
6406     \flag_if_raised:nT { str_end }
6407     {
6408         \\
6409         The-string-contained-an-odd-number-of-bytes.-This-is-invalid:~
6410         the-basic-code-unit-for-UTF-16-is-16-bits-(2-bytes).
6411     }
6412 }

```

(End definition for `\l__str_missing_flag`, `\l__str_extra_flag`, and `\l__str_end_flag`.)

`__str_convert_decode_utf16:` As for UTF-8, decoding UTF-16 is harder than encoding it. If the endianness is unknown, check the first two bytes: if those are "FE and "FF in either order, remove them and use the corresponding endianness, otherwise assume big-endianness. The three endianness cases are based on a common auxiliary whose first argument is 1 for big-endian and 2 for little-endian, and whose second argument, delimited by the scan mark `\s__str_stop`, is expanded once (the string may be long; passing `\g__str_result_tl` as an argument before expansion is cheaper).

The `__str_decode_utf_xvi:Nw` function defines `__str_tmp:w` to take two arguments and return the character code of the first one if the string is big-endian, and the second one if the string is little-endian, then loops over the string using `__str_decode_utf_xvi_pair:NN` described below.

```

6413     \cs_new_protected:cpn { __str_convert_decode_utf16be: }
6414     { \__str_decode_utf_xvi:Nw 1 \g__str_result_tl \s__str_stop }
6415     \cs_new_protected:cpn { __str_convert_decode_utf16le: }
6416     { \__str_decode_utf_xvi:Nw 2 \g__str_result_tl \s__str_stop }
6417     \cs_new_protected:cpn { __str_convert_decode_utf16: }
6418     {
6419         \exp_after:wN \__str_decode_utf_xvi_bom:NN
6420         \g__str_result_tl \s__str_stop \s__str_stop \s__str_stop
6421     }
6422     \cs_new_protected:Npn \__str_decode_utf_xvi_bom:NN #1#2
6423     {
6424         \str_if_eq:nnTF { #1#2 } { ^^ff ^^fe }
6425         { \__str_decode_utf_xvi:Nw 2 }
6426         {
6427             \str_if_eq:nnTF { #1#2 } { ^^fe ^^ff }
6428             { \__str_decode_utf_xvi:Nw 1 }
6429             { \__str_decode_utf_xvi:Nw 1 #1#2 }
6430         }

```

```

6431     }
6432     \cs_new_protected:Npn \__str_decode_utf_xvi:Nw #1#2 \s__str_stop
6433     {
6434         \flag_clear:n { str_error }
6435         \flag_clear:n { str_missing }
6436         \flag_clear:n { str_extra }
6437         \flag_clear:n { str_end }
6438         \cs_set:Npn \__str_tmp:w ##1 ##2 { ' ## #1 }
6439         \__kernel_tl_gset:Nx \g__str_result_tl
6440         {
6441             \exp_after:wN \__str_decode_utf_xvi_pair:NN
6442             #2 \q__str_nil \q__str_nil
6443             \prg_break_point:
6444         }
6445         \__str_if_flag_error:nmx { str_error } { utf16-decode } { }
6446     }

```

(End definition for `__str_convert_decode_utf16:` and others.)

```

\__str_decode_utf_xvi_pair:NN
\__str_decode_utf_xvi_quad:NNwNN
\__str_decode_utf_xvi_pair_end:Nw
\__str_decode_utf_xvi_error:nNN
\__str_decode_utf_xvi_extra:NNw

```

Bytes are read two at a time. At this stage, `\@@_tmp:w #1#2` expands to the character code of the most significant byte, and we distinguish cases depending on which range it lies in:

- ["D8, "DB] signals a lead surrogate, and the integer expression yields 1 (ε -TeX rounds ties away from zero);
- ["DC, "DF] signals a trail surrogate, unexpected here, and the integer expression yields 2;
- any other value signals a code point in the Basic Multilingual Plane, which stands for itself, and the `\if_case:w` construction expands to nothing (cases other than 1 or 2), leaving the relevant material in the input stream, followed by another call to the `_pair` auxiliary.

The case of a lead surrogate is treated by the `_quad` auxiliary, whose arguments `#1`, `#2`, `#4` and `#5` are the four bytes. We expect the most significant byte of `#4#5` to be in the range ["DC, "DF] (trail surrogate). The test is similar to the test used for continuation bytes in the UTF-8 decoding functions. In the case where `#4#5` is indeed a trail surrogate, leave `#1#2#4#5 \s__str <code point> \s__str`, and remove the pair `#4#5` before looping with `__str_decode_utf_xvi_pair:NN`. Otherwise, of course, complain about the missing surrogate.

The magic number "D7F7 is such that $"D7F7*"400 = "D800*"400 + "DC00 - "10000$.

Every time we read a pair of bytes, we test for the end-marker `\q__str_nil`. When reaching the end, we additionally check that the string had an even length. Also, if the end is reached when expecting a trail surrogate, we treat that as a missing surrogate.

```

6447     \cs_new:Npn \__str_decode_utf_xvi_pair:NN #1#2
6448     {
6449         \if_meaning:w \q__str_nil #2
6450         \__str_decode_utf_xvi_pair_end:Nw #1
6451         \fi:
6452         \if_case:w
6453             \int_eval:n { ( \__str_tmp:w #1#2 - "D6 ) / 4 } \scan_stop:
6454         \or: \exp_after:wN \__str_decode_utf_xvi_quad:NNwNN
6455         \or: \exp_after:wN \__str_decode_utf_xvi_extra:NNw

```

```

6456     \fi:
6457     #1#2 \s__str
6458     \int_eval:n { "100 * \__str_tmp:w #1#2 + \__str_tmp:w #2#1 } \s__str
6459     \__str_decode_utf_xvi_pair:NN
6460 }
6461 \cs_new:Npn \__str_decode_utf_xvi_quad:NNwNN
6462   #1#2 #3 \__str_decode_utf_xvi_pair:NN #4#5
6463 {
6464   \if_meaning:w \q__str_nil #5
6465     \__str_decode_utf_xvi_error:nNN { missing } #1#2
6466     \__str_decode_utf_xvi_pair_end:Nw #4
6467   \fi:
6468   \if_int_compare:w
6469     \if_int_compare:w \__str_tmp:w #4#5 < "DC \exp_stop_f:
6470       0 = 1
6471     \else:
6472       \__str_tmp:w #4#5 < "E0
6473     \fi:
6474     \exp_stop_f:
6475     #1 #2 #4 #5 \s__str
6476     \int_eval:n
6477     {
6478       ( "100 * \__str_tmp:w #1#2 + \__str_tmp:w #2#1 - "D7F7 ) * "400
6479       + "100 * \__str_tmp:w #4#5 + \__str_tmp:w #5#4
6480     }
6481     \s__str
6482     \exp_after:wN \use_i:nnn
6483   \else:
6484     \__str_decode_utf_xvi_error:nNN { missing } #1#2
6485   \fi:
6486   \__str_decode_utf_xvi_pair:NN #4#5
6487 }
6488 \cs_new:Npn \__str_decode_utf_xvi_pair_end:Nw #1 \fi:
6489 {
6490   \fi:
6491   \if_meaning:w \q__str_nil #1
6492   \else:
6493     \__str_decode_utf_xvi_error:nNN { end } #1 \prg_do_nothing:
6494   \fi:
6495   \prg_break:
6496 }
6497 \cs_new:Npn \__str_decode_utf_xvi_extra:NNw #1#2 \s__str #3 \s__str
6498 { \__str_decode_utf_xvi_error:nNN { extra } #1#2 }
6499 \cs_new:Npn \__str_decode_utf_xvi_error:nNN #1#2#3
6500 {
6501   \flag_raise:n { str_error }
6502   \flag_raise:n { str_#1 }
6503   #2 #3 \s__str
6504   \int_use:N \c__str_replacement_char_int \s__str
6505 }

```

(End definition for __str_decode_utf_xvi_pair:NN and others.)

Restore the original catcodes of bytes 254 and 255.

```

6506 \group_end:

```

10.6.3 utf-32 support

The definitions are done in a category code regime where the bytes 0, 254 and 255 used by the byte order mark have catcode “other”.

```
6507 \group_begin:
6508   \char_set_catcode_other:N \^^00
6509   \char_set_catcode_other:N \^^fe
6510   \char_set_catcode_other:N \^^ff
```

`__str_convert_encode_utf32:` Convert each integer in the comma-list `\g__str_result_tl` to a sequence of four bytes. The functions for big-endian and little-endian encodings are very similar, but the `__str_output_byte:n` instructions are reversed.

```
\__str_convert_encode_utf32be:
  \__str_convert_encode_utf32le:
\__str_encode_utf_xxxii_be:n
  \__str_encode_utf_xxxii_be_aux:nn
\__str_encode_utf_xxxii_le:n
  \__str_encode_utf_xxxii_le_aux:nn
6511   \cs_new_protected:cpn { __str_convert_encode_utf32: }
6512   {
6513     \__str_convert_gmap_internal:N \__str_encode_utf_xxxii_be:n
6514     \tl_gput_left:Nx \g__str_result_tl { ^^00 ^^00 ^^fe ^^ff }
6515   }
6516   \cs_new_protected:cpn { __str_convert_encode_utf32be: }
6517   { \__str_convert_gmap_internal:N \__str_encode_utf_xxxii_be:n }
6518   \cs_new_protected:cpn { __str_convert_encode_utf32le: }
6519   { \__str_convert_gmap_internal:N \__str_encode_utf_xxxii_le:n }
6520   \cs_new:Npn \__str_encode_utf_xxxii_be:n #1
6521   {
6522     \exp_args:Nf \__str_encode_utf_xxxii_be_aux:nn
6523     { \int_div_truncate:nn {#1} { "100 } } {#1}
6524   }
6525   \cs_new:Npn \__str_encode_utf_xxxii_be_aux:nn #1#2
6526   {
6527     ^^00
6528     \__str_output_byte_pair_be:n {#1}
6529     \__str_output_byte:n { #2 - #1 * "100 }
6530   }
6531   \cs_new:Npn \__str_encode_utf_xxxii_le:n #1
6532   {
6533     \exp_args:Nf \__str_encode_utf_xxxii_le_aux:nn
6534     { \int_div_truncate:nn {#1} { "100 } } {#1}
6535   }
6536   \cs_new:Npn \__str_encode_utf_xxxii_le_aux:nn #1#2
6537   {
6538     \__str_output_byte:n { #2 - #1 * "100 }
6539     \__str_output_byte_pair_le:n {#1}
6540     ^^00
6541   }
```

(End definition for `__str_convert_encode_utf32:` and others.)

str_overflow There can be no error when encoding in UTF-32. When decoding, the string may not have length $4n$, or it may contain code points larger than "10FFFF". The latter case often happens if the encoding was in fact not UTF-32, because most arbitrary strings are not valid in UTF-32.

```
6542   \flag_clear_new:n { str_overflow }
6543   \flag_clear_new:n { str_end }
6544   \__kernel_msg_new:nnnn { str } { utf32-decode }
6545   {
```

```

6546 Invalid-UTF-32-string:
6547 \exp_last_unbraced:Nf \use_none:n
6548 {
6549   \__str_if_flag_times:nT { str_overflow } { ,~code-point-too-large }
6550   \__str_if_flag_times:nT { str_end }      { ,~truncated-string }
6551 }
6552 .
6553 }
6554 {
6555   In~the~UTF-32~encoding,~every~Unicode~character~
6556   (in~the~range~[U+0000,~U+10FFFF])~is~encoded~as~4~bytes.
6557   \flag_if_raised:nT { str_overflow }
6558   {
6559     \\\
6560     LaTeX~came~across~a~code~point~larger~than~1114111,~
6561     the~maximum~code~point~defined~by~Unicode.~
6562     Perhaps~the~string~was~not~encoded~in~the~UTF-32~encoding?
6563   }
6564   \flag_if_raised:nT { str_end }
6565   {
6566     \\\
6567     The~length~of~the~string~is~not~a~multiple~of~4.~
6568     Perhaps~the~string~was~truncated?
6569   }
6570 }

```

(End definition for `str_overflow` and `str_end`. These variables are documented on page ??.)

`__str_convert_decode_utf32:` The structure is similar to UTF-16 decoding functions. If the endianness is not given, test the first 4 bytes of the string (possibly `\s__str_stop` if the string is too short) for the presence of a byte-order mark. If there is a byte-order mark, use that endianness, and remove the 4 bytes, otherwise default to big-endian, and leave the 4 bytes in place. The `__str_decode_utf_xxxii:Nw` auxiliary receives 1 or 2 as its first argument indicating endianness, and the string to convert as its second argument (expanded or not). It sets `__str_tmp:w` to expand to the character code of either of its two arguments depending on endianness, then triggers the `_loop` auxiliary inside an x-expanding assignment to `\g__str_result_tl`.

The `_loop` auxiliary first checks for the end-of-string marker `\s__str_stop`, calling the `_end` auxiliary if appropriate. Otherwise, leave the $\langle 4 \text{ bytes} \rangle$ `\s__str` behind, then check that the code point is not overflowing: the leading byte must be 0, and the following byte at most 16.

In the ending code, we check that there remains no byte: there should be nothing left until the first `\s__str_stop`. Break the map.

```

6571 \cs_new_protected:cpn { __str_convert_decode_utf32be: }
6572 { \__str_decode_utf_xxxii:Nw 1 \g__str_result_tl \s__str_stop }
6573 \cs_new_protected:cpn { __str_convert_decode_utf32le: }
6574 { \__str_decode_utf_xxxii:Nw 2 \g__str_result_tl \s__str_stop }
6575 \cs_new_protected:cpn { __str_convert_decode_utf32: }
6576 {
6577   \exp_after:wN \__str_decode_utf_xxxii_bom:NNNN \g__str_result_tl
6578   \s__str_stop \s__str_stop \s__str_stop \s__str_stop \s__str_stop
6579 }
6580 \cs_new_protected:Npn \__str_decode_utf_xxxii_bom:NNNN #1#2#3#4

```

```

6581 {
6582     \str_if_eq:nnTF { #1#2#3#4 } { ^^ff ^^fe ^^00 ^^00 }
6583     { \__str_decode_utf_xxxii:Nw 2 }
6584     {
6585         \str_if_eq:nnTF { #1#2#3#4 } { ^^00 ^^00 ^^fe ^^ff }
6586         { \__str_decode_utf_xxxii:Nw 1 }
6587         { \__str_decode_utf_xxxii:Nw 1 #1#2#3#4 }
6588     }
6589 }
6590 \cs_new_protected:Npn \__str_decode_utf_xxxii:Nw #1#2 \s__str_stop
6591 {
6592     \flag_clear:n { str_overflow }
6593     \flag_clear:n { str_end }
6594     \flag_clear:n { str_error }
6595     \cs_set:Npn \__str_tmp:w ##1 ##2 { ' ## #1 }
6596     \__kernel_tl_gset:Nx \g__str_result_tl
6597     {
6598         \exp_after:wN \__str_decode_utf_xxxii_loop:NNNN
6599         #2 \s__str_stop \s__str_stop \s__str_stop \s__str_stop
6600         \prg_break_point:
6601     }
6602     \__str_if_flag_error:nmx { str_error } { utf32-decode } { }
6603 }
6604 \cs_new:Npn \__str_decode_utf_xxxii_loop:NNNN #1#2#3#4
6605 {
6606     \if_meaning:w \s__str_stop #4
6607     \exp_after:wN \__str_decode_utf_xxxii_end:w
6608     \fi:
6609     #1#2#3#4 \s__str
6610     \if_int_compare:w \__str_tmp:w #1#4 > 0 \exp_stop_f:
6611         \flag_raise:n { str_overflow }
6612         \flag_raise:n { str_error }
6613         \int_use:N \c__str_replacement_char_int
6614     \else:
6615         \if_int_compare:w \__str_tmp:w #2#3 > 16 \exp_stop_f:
6616             \flag_raise:n { str_overflow }
6617             \flag_raise:n { str_error }
6618             \int_use:N \c__str_replacement_char_int
6619         \else:
6620             \int_eval:n
6621             { \__str_tmp:w #2#3*"10000 + \__str_tmp:w #3#2*"100 + \__str_tmp:w #4#1 }
6622             \fi:
6623         \fi:
6624         \s__str
6625         \__str_decode_utf_xxxii_loop:NNNN
6626     }
6627 \cs_new:Npn \__str_decode_utf_xxxii_end:w #1 \s__str_stop
6628 {
6629     \tl_if_empty:nF {#1}
6630     {
6631         \flag_raise:n { str_end }
6632         \flag_raise:n { str_error }
6633         #1 \s__str
6634         \int_use:N \c__str_replacement_char_int \s__str

```

```

6635     }
6636     \prg_break:
6637 }

```

(End definition for `_str_convert_decode_utf32:` and others.)

Restore the original catcodes of bytes 0, 254 and 255.

```

6638 \group_end:

```

10.7 PDF names and strings by expansion

```

\str_convert_pdfname:n
\_str_convert_pdfname:n
\_str_convert_pdfname_bytes:n
\_str_convert_pdfname_bytes_aux:n
\_str_convert_pdfname_bytes_aux:nnn

```

To convert to PDF names by expansion, we work purely on UTF-8 input. The first step is to make a string with “other” spaces, after which we use a simple token-by-token approach. In Unicode engines, we break down everything before one-byte codepoints, but for 8-bit engines there is no need to worry. Actual escaping is covered by the same code as used in the non-expandable route.

```

6639 \cs_new:Npn \str_convert_pdfname:n #1
6640 {
6641   \exp_args:Ne \tl_to_str:n
6642   { \str_map_function:nN {#1} \_str_convert_pdfname:n }
6643 }
6644 \bool_lazy_or:nnTF
6645 { \sys_if_engine luatex_p: }
6646 { \sys_if_engine xetex_p: }
6647 {
6648   \cs_new:Npn \_str_convert_pdfname:n #1
6649   {
6650     \int_compare:nNnTF { '#1 } > { "7F }
6651     { \_str_convert_pdfname_bytes:n {#1} }
6652     { \_str_escape_name_char:n {#1} }
6653   }
6654   \cs_new:Npn \_str_convert_pdfname_bytes:n #1
6655   {
6656     \exp_args:Ne \_str_convert_pdfname_bytes_aux:n
6657     { \char_to_utfviii_bytes:n {'#1} }
6658   }
6659   \cs_new:Npn \_str_convert_pdfname_bytes_aux:n #1
6660   { \_str_convert_pdfname_bytes_aux:nnnn #1 }
6661   \cs_new:Npx \_str_convert_pdfname_bytes_aux:nnnn #1#2#3#4
6662   {
6663     \c_hash_str \exp_not:N \_str_output_hexadecimal:n {#1}
6664     \c_hash_str \exp_not:N \_str_output_hexadecimal:n {#2}
6665     \exp_not:N \tl_if_blank:nF {#3}
6666     {
6667       \c_hash_str \exp_not:N \_str_output_hexadecimal:n {#3}
6668       \exp_not:N \tl_if_blank:nF {#4}
6669       {
6670         \c_hash_str \exp_not:N \_str_output_hexadecimal:n {#4}
6671       }
6672     }
6673   }
6674 }
6675 { \cs_new_eq:NN \_str_convert_pdfname:n \_str_escape_name_char:n }

```

(End definition for `\str_convert_pdfname:n` and others. This function is documented on page 74.)

```
6676 </package>
```

10.7.1 iso 8859 support

The ISO-8859-1 encoding exactly matches with the 256 first Unicode characters. For other 8-bit encodings of the ISO-8859 family, we keep track only of differences, and of unassigned bytes.

```
6677 <*iso88591>
6678 \__str_declare_eight_bit_encoding:nnnn { iso88591 } { 256 }
6679 {
6680 }
6681 {
6682 }
6683 </iso88591>

6684 <*iso88592>
6685 \__str_declare_eight_bit_encoding:nnnn { iso88592 } { 399 }
6686 {
6687   { A1 } { 0104 }
6688   { A2 } { 02D8 }
6689   { A3 } { 0141 }
6690   { A5 } { 013D }
6691   { A6 } { 015A }
6692   { A9 } { 0160 }
6693   { AA } { 015E }
6694   { AB } { 0164 }
6695   { AC } { 0179 }
6696   { AE } { 017D }
6697   { AF } { 017B }
6698   { B1 } { 0105 }
6699   { B2 } { 02DB }
6700   { B3 } { 0142 }
6701   { B5 } { 013E }
6702   { B6 } { 015B }
6703   { B7 } { 02C7 }
6704   { B9 } { 0161 }
6705   { BA } { 015F }
6706   { BB } { 0165 }
6707   { BC } { 017A }
6708   { BD } { 02DD }
6709   { BE } { 017E }
6710   { BF } { 017C }
6711   { C0 } { 0154 }
6712   { C3 } { 0102 }
6713   { C5 } { 0139 }
6714   { C6 } { 0106 }
6715   { C8 } { 010C }
6716   { CA } { 0118 }
6717   { CC } { 011A }
6718   { CF } { 010E }
6719   { D0 } { 0110 }
6720   { D1 } { 0143 }
6721   { D2 } { 0147 }
```



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6722     { D5 } { 0150 }
6723     { D8 } { 0158 }
6724     { D9 } { 016E }
6725     { DB } { 0170 }
6726     { DE } { 0162 }
6727     { E0 } { 0155 }
6728     { E3 } { 0103 }
6729     { E5 } { 013A }
6730     { E6 } { 0107 }
6731     { E8 } { 010D }
6732     { EA } { 0119 }
6733     { EC } { 011B }
6734     { EF } { 010F }
6735     { F0 } { 0111 }
6736     { F1 } { 0144 }
6737     { F2 } { 0148 }
6738     { F5 } { 0151 }
6739     { F8 } { 0159 }
6740     { F9 } { 016F }
6741     { FB } { 0171 }
6742     { FE } { 0163 }
6743     { FF } { 02D9 }
6744 }
6745 {
6746 }
6747 </iso88592>
6748 <iso88593>
6749 \__str_declare_eight_bit_encoding:nmmn { iso88593 } { 384 }
6750 {
6751     { A1 } { 0126 }
6752     { A2 } { 02D8 }
6753     { A6 } { 0124 }
6754     { A9 } { 0130 }
6755     { AA } { 015E }
6756     { AB } { 011E }
6757     { AC } { 0134 }
6758     { AF } { 017B }
6759     { B1 } { 0127 }
6760     { B6 } { 0125 }
6761     { B9 } { 0131 }
6762     { BA } { 015F }
6763     { BB } { 011F }
6764     { BC } { 0135 }
6765     { BF } { 017C }
6766     { C5 } { 010A }
6767     { C6 } { 0108 }
6768     { D5 } { 0120 }
6769     { D8 } { 011C }
6770     { DD } { 016C }
6771     { DE } { 015C }
6772     { E5 } { 010B }
6773     { E6 } { 0109 }
6774     { F5 } { 0121 }
6775     { F8 } { 011D }

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6776     { FD } { 016D }
6777     { FE } { 015D }
6778     { FF } { 02D9 }
6779 }
6780 {
6781     { A5 }
6782     { AE }
6783     { BE }
6784     { C3 }
6785     { D0 }
6786     { E3 }
6787     { F0 }
6788 }
6789 </iso88593>
6790 <*iso88594>
6791 \__str_declareEight_bit_encoding:nnnn { iso88594 } { 383 }
6792 {
6793     { A1 } { 0104 }
6794     { A2 } { 0138 }
6795     { A3 } { 0156 }
6796     { A5 } { 0128 }
6797     { A6 } { 013B }
6798     { A9 } { 0160 }
6799     { AA } { 0112 }
6800     { AB } { 0122 }
6801     { AC } { 0166 }
6802     { AE } { 017D }
6803     { B1 } { 0105 }
6804     { B2 } { 02DB }
6805     { B3 } { 0157 }
6806     { B5 } { 0129 }
6807     { B6 } { 013C }
6808     { B7 } { 02C7 }
6809     { B9 } { 0161 }
6810     { BA } { 0113 }
6811     { BB } { 0123 }
6812     { BC } { 0167 }
6813     { BD } { 014A }
6814     { BE } { 017E }
6815     { BF } { 014B }
6816     { C0 } { 0100 }
6817     { C7 } { 012E }
6818     { C8 } { 010C }
6819     { CA } { 0118 }
6820     { CC } { 0116 }
6821     { CF } { 012A }
6822     { D0 } { 0110 }
6823     { D1 } { 0145 }
6824     { D2 } { 014C }
6825     { D3 } { 0136 }
6826     { D9 } { 0172 }
6827     { DD } { 0168 }
6828     { DE } { 016A }
6829     { EO } { 0101 }

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6830     { E7 } { 012F }
6831     { E8 } { 010D }
6832     { EA } { 0119 }
6833     { EC } { 0117 }
6834     { EF } { 012B }
6835     { F0 } { 0111 }
6836     { F1 } { 0146 }
6837     { F2 } { 014D }
6838     { F3 } { 0137 }
6839     { F9 } { 0173 }
6840     { FD } { 0169 }
6841     { FE } { 016B }
6842     { FF } { 02D9 }
6843 }
6844 {
6845 }
6846 </iso88594>
6847 <*:iso88595>
6848 \__str_declare_eight_bit_encoding:nnnn { iso88595 } { 374 }
6849 {
6850     { A1 } { 0401 }
6851     { A2 } { 0402 }
6852     { A3 } { 0403 }
6853     { A4 } { 0404 }
6854     { A5 } { 0405 }
6855     { A6 } { 0406 }
6856     { A7 } { 0407 }
6857     { A8 } { 0408 }
6858     { A9 } { 0409 }
6859     { AA } { 040A }
6860     { AB } { 040B }
6861     { AC } { 040C }
6862     { AE } { 040E }
6863     { AF } { 040F }
6864     { B0 } { 0410 }
6865     { B1 } { 0411 }
6866     { B2 } { 0412 }
6867     { B3 } { 0413 }
6868     { B4 } { 0414 }
6869     { B5 } { 0415 }
6870     { B6 } { 0416 }
6871     { B7 } { 0417 }
6872     { B8 } { 0418 }
6873     { B9 } { 0419 }
6874     { BA } { 041A }
6875     { BB } { 041B }
6876     { BC } { 041C }
6877     { BD } { 041D }
6878     { BE } { 041E }
6879     { BF } { 041F }
6880     { C0 } { 0420 }
6881     { C1 } { 0421 }
6882     { C2 } { 0422 }
6883     { C3 } { 0423 }

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6884	{ C4 }	{ 0424 }
6885	{ C5 }	{ 0425 }
6886	{ C6 }	{ 0426 }
6887	{ C7 }	{ 0427 }
6888	{ C8 }	{ 0428 }
6889	{ C9 }	{ 0429 }
6890	{ CA }	{ 042A }
6891	{ CB }	{ 042B }
6892	{ CC }	{ 042C }
6893	{ CD }	{ 042D }
6894	{ CE }	{ 042E }
6895	{ CF }	{ 042F }
6896	{ D0 }	{ 0430 }
6897	{ D1 }	{ 0431 }
6898	{ D2 }	{ 0432 }
6899	{ D3 }	{ 0433 }
6900	{ D4 }	{ 0434 }
6901	{ D5 }	{ 0435 }
6902	{ D6 }	{ 0436 }
6903	{ D7 }	{ 0437 }
6904	{ D8 }	{ 0438 }
6905	{ D9 }	{ 0439 }
6906	{ DA }	{ 043A }
6907	{ DB }	{ 043B }
6908	{ DC }	{ 043C }
6909	{ DD }	{ 043D }
6910	{ DE }	{ 043E }
6911	{ DF }	{ 043F }
6912	{ E0 }	{ 0440 }
6913	{ E1 }	{ 0441 }
6914	{ E2 }	{ 0442 }
6915	{ E3 }	{ 0443 }
6916	{ E4 }	{ 0444 }
6917	{ E5 }	{ 0445 }
6918	{ E6 }	{ 0446 }
6919	{ E7 }	{ 0447 }
6920	{ E8 }	{ 0448 }
6921	{ E9 }	{ 0449 }
6922	{ EA }	{ 044A }
6923	{ EB }	{ 044B }
6924	{ EC }	{ 044C }
6925	{ ED }	{ 044D }
6926	{ EE }	{ 044E }
6927	{ EF }	{ 044F }
6928	{ F0 }	{ 2116 }
6929	{ F1 }	{ 0451 }
6930	{ F2 }	{ 0452 }
6931	{ F3 }	{ 0453 }
6932	{ F4 }	{ 0454 }
6933	{ F5 }	{ 0455 }
6934	{ F6 }	{ 0456 }
6935	{ F7 }	{ 0457 }
6936	{ F8 }	{ 0458 }
6937	{ F9 }	{ 0459 }

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6938     { FA } { 045A }
6939     { FB } { 045B }
6940     { FC } { 045C }
6941     { FD } { 00A7 }
6942     { FE } { 045E }
6943     { FF } { 045F }
6944 }
6945 {
6946 }
6947 </iso88595>
6948 <*iso88596>
6949 \_str_declare\_eight\_bit\_encoding:nnnn { iso88596 } { 344 }
6950 {
6951     { AC } { 060C }
6952     { BB } { 061B }
6953     { BF } { 061F }
6954     { C1 } { 0621 }
6955     { C2 } { 0622 }
6956     { C3 } { 0623 }
6957     { C4 } { 0624 }
6958     { C5 } { 0625 }
6959     { C6 } { 0626 }
6960     { C7 } { 0627 }
6961     { C8 } { 0628 }
6962     { C9 } { 0629 }
6963     { CA } { 062A }
6964     { CB } { 062B }
6965     { CC } { 062C }
6966     { CD } { 062D }
6967     { CE } { 062E }
6968     { CF } { 062F }
6969     { D0 } { 0630 }
6970     { D1 } { 0631 }
6971     { D2 } { 0632 }
6972     { D3 } { 0633 }
6973     { D4 } { 0634 }
6974     { D5 } { 0635 }
6975     { D6 } { 0636 }
6976     { D7 } { 0637 }
6977     { D8 } { 0638 }
6978     { D9 } { 0639 }
6979     { DA } { 063A }
6980     { E0 } { 0640 }
6981     { E1 } { 0641 }
6982     { E2 } { 0642 }
6983     { E3 } { 0643 }
6984     { E4 } { 0644 }
6985     { E5 } { 0645 }
6986     { E6 } { 0646 }
6987     { E7 } { 0647 }
6988     { E8 } { 0648 }
6989     { E9 } { 0649 }
6990     { EA } { 064A }
6991     { EB } { 064B }

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6992     { EC } { 064C }
6993     { ED } { 064D }
6994     { EE } { 064E }
6995     { EF } { 064F }
6996     { F0 } { 0650 }
6997     { F1 } { 0651 }
6998     { F2 } { 0652 }
6999 }
7000 {
7001     { A1 }
7002     { A2 }
7003     { A3 }
7004     { A5 }
7005     { A6 }
7006     { A7 }
7007     { A8 }
7008     { A9 }
7009     { AA }
7010     { AB }
7011     { AE }
7012     { AF }
7013     { B0 }
7014     { B1 }
7015     { B2 }
7016     { B3 }
7017     { B4 }
7018     { B5 }
7019     { B6 }
7020     { B7 }
7021     { B8 }
7022     { B9 }
7023     { BA }
7024     { BC }
7025     { BD }
7026     { BE }
7027     { C0 }
7028     { DB }
7029     { DC }
7030     { DD }
7031     { DE }
7032     { DF }
7033 }
7034 </iso88596>
7035 <iso88597>
7036 \__str_declare_eight_bit_encoding:nmn { iso88597 } { 498 }
7037 {
7038     { A1 } { 2018 }
7039     { A2 } { 2019 }
7040     { A4 } { 20AC }
7041     { A5 } { 20AF }
7042     { AA } { 037A }
7043     { AF } { 2015 }
7044     { B4 } { 0384 }
7045     { B5 } { 0385 }

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7046	{ B6 }	{ 0386 }
7047	{ B8 }	{ 0388 }
7048	{ B9 }	{ 0389 }
7049	{ BA }	{ 038A }
7050	{ BC }	{ 038C }
7051	{ BE }	{ 038E }
7052	{ BF }	{ 038F }
7053	{ C0 }	{ 0390 }
7054	{ C1 }	{ 0391 }
7055	{ C2 }	{ 0392 }
7056	{ C3 }	{ 0393 }
7057	{ C4 }	{ 0394 }
7058	{ C5 }	{ 0395 }
7059	{ C6 }	{ 0396 }
7060	{ C7 }	{ 0397 }
7061	{ C8 }	{ 0398 }
7062	{ C9 }	{ 0399 }
7063	{ CA }	{ 039A }
7064	{ CB }	{ 039B }
7065	{ CC }	{ 039C }
7066	{ CD }	{ 039D }
7067	{ CE }	{ 039E }
7068	{ CF }	{ 039F }
7069	{ D0 }	{ 03A0 }
7070	{ D1 }	{ 03A1 }
7071	{ D3 }	{ 03A3 }
7072	{ D4 }	{ 03A4 }
7073	{ D5 }	{ 03A5 }
7074	{ D6 }	{ 03A6 }
7075	{ D7 }	{ 03A7 }
7076	{ D8 }	{ 03A8 }
7077	{ D9 }	{ 03A9 }
7078	{ DA }	{ 03AA }
7079	{ DB }	{ 03AB }
7080	{ DC }	{ 03AC }
7081	{ DD }	{ 03AD }
7082	{ DE }	{ 03AE }
7083	{ DF }	{ 03AF }
7084	{ E0 }	{ 03B0 }
7085	{ E1 }	{ 03B1 }
7086	{ E2 }	{ 03B2 }
7087	{ E3 }	{ 03B3 }
7088	{ E4 }	{ 03B4 }
7089	{ E5 }	{ 03B5 }
7090	{ E6 }	{ 03B6 }
7091	{ E7 }	{ 03B7 }
7092	{ E8 }	{ 03B8 }
7093	{ E9 }	{ 03B9 }
7094	{ EA }	{ 03BA }
7095	{ EB }	{ 03BB }
7096	{ EC }	{ 03BC }
7097	{ ED }	{ 03BD }
7098	{ EE }	{ 03BE }
7099	{ EF }	{ 03BF }

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7100     { F0 } { 03C0 }
7101     { F1 } { 03C1 }
7102     { F2 } { 03C2 }
7103     { F3 } { 03C3 }
7104     { F4 } { 03C4 }
7105     { F5 } { 03C5 }
7106     { F6 } { 03C6 }
7107     { F7 } { 03C7 }
7108     { F8 } { 03C8 }
7109     { F9 } { 03C9 }
7110     { FA } { 03CA }
7111     { FB } { 03CB }
7112     { FC } { 03CC }
7113     { FD } { 03CD }
7114     { FE } { 03CE }
7115 }
7116 {
7117     { AE }
7118     { D2 }
7119 }
7120 </iso88597>
7121 < *iso88598>
7122 \_str_declare\_eight\_bit\_encoding:nnnn { iso88598 } { 308 }
7123 {
7124     { AA } { 00D7 }
7125     { BA } { 00F7 }
7126     { DF } { 2017 }
7127     { E0 } { 05D0 }
7128     { E1 } { 05D1 }
7129     { E2 } { 05D2 }
7130     { E3 } { 05D3 }
7131     { E4 } { 05D4 }
7132     { E5 } { 05D5 }
7133     { E6 } { 05D6 }
7134     { E7 } { 05D7 }
7135     { E8 } { 05D8 }
7136     { E9 } { 05D9 }
7137     { EA } { 05DA }
7138     { EB } { 05DB }
7139     { EC } { 05DC }
7140     { ED } { 05DD }
7141     { EE } { 05DE }
7142     { EF } { 05DF }
7143     { F0 } { 05E0 }
7144     { F1 } { 05E1 }
7145     { F2 } { 05E2 }
7146     { F3 } { 05E3 }
7147     { F4 } { 05E4 }
7148     { F5 } { 05E5 }
7149     { F6 } { 05E6 }
7150     { F7 } { 05E7 }
7151     { F8 } { 05E8 }
7152     { F9 } { 05E9 }
7153     { FA } { 05EA }

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7154     { FD } { 200E }
7155     { FE } { 200F }
7156 }
7157 {
7158     { A1 }
7159     { BF }
7160     { C0 }
7161     { C1 }
7162     { C2 }
7163     { C3 }
7164     { C4 }
7165     { C5 }
7166     { C6 }
7167     { C7 }
7168     { C8 }
7169     { C9 }
7170     { CA }
7171     { CB }
7172     { CC }
7173     { CD }
7174     { CE }
7175     { CF }
7176     { D0 }
7177     { D1 }
7178     { D2 }
7179     { D3 }
7180     { D4 }
7181     { D5 }
7182     { D6 }
7183     { D7 }
7184     { D8 }
7185     { D9 }
7186     { DA }
7187     { DB }
7188     { DC }
7189     { DD }
7190     { DE }
7191     { FB }
7192     { FC }
7193 }
7194 </iso88598>
7195 < *iso88599>
7196 \_str_declare\_eight\_bit\_encoding:nnnn { iso88599 } { 352 }
7197 {
7198     { D0 } { 011E }
7199     { DD } { 0130 }
7200     { DE } { 015E }
7201     { FO } { 011F }
7202     { FD } { 0131 }
7203     { FE } { 015F }
7204 }
7205 {
7206 }
7207 </iso88599>

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```

7208 <*iso885910>
7209 \_str_declare_eight_bit_encoding:nnnn { iso885910 } { 383 }
7210 {
7211     { A1 } { 0104 }
7212     { A2 } { 0112 }
7213     { A3 } { 0122 }
7214     { A4 } { 012A }
7215     { A5 } { 0128 }
7216     { A6 } { 0136 }
7217     { A8 } { 013B }
7218     { A9 } { 0110 }
7219     { AA } { 0160 }
7220     { AB } { 0166 }
7221     { AC } { 017D }
7222     { AE } { 016A }
7223     { AF } { 014A }
7224     { B1 } { 0105 }
7225     { B2 } { 0113 }
7226     { B3 } { 0123 }
7227     { B4 } { 012B }
7228     { B5 } { 0129 }
7229     { B6 } { 0137 }
7230     { B8 } { 013C }
7231     { B9 } { 0111 }
7232     { BA } { 0161 }
7233     { BB } { 0167 }
7234     { BC } { 017E }
7235     { BD } { 2015 }
7236     { BE } { 016B }
7237     { BF } { 014B }
7238     { C0 } { 0100 }
7239     { C7 } { 012E }
7240     { C8 } { 010C }
7241     { CA } { 0118 }
7242     { CC } { 0116 }
7243     { D1 } { 0145 }
7244     { D2 } { 014C }
7245     { D7 } { 0168 }
7246     { D9 } { 0172 }
7247     { E0 } { 0101 }
7248     { E7 } { 012F }
7249     { E8 } { 010D }
7250     { EA } { 0119 }
7251     { EC } { 0117 }
7252     { F1 } { 0146 }
7253     { F2 } { 014D }
7254     { F7 } { 0169 }
7255     { F9 } { 0173 }
7256     { FF } { 0138 }
7257 }
7258 {
7259 }
7260 </iso885910>
7261 <*iso885911>

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```

7262 \__str_declare_eight_bit_encoding:nnnn { iso885911 } { 369 }
7263 {
7264     { A1 } { OE01 }
7265     { A2 } { OE02 }
7266     { A3 } { OE03 }
7267     { A4 } { OE04 }
7268     { A5 } { OE05 }
7269     { A6 } { OE06 }
7270     { A7 } { OE07 }
7271     { A8 } { OE08 }
7272     { A9 } { OE09 }
7273     { AA } { OE0A }
7274     { AB } { OE0B }
7275     { AC } { OE0C }
7276     { AD } { OE0D }
7277     { AE } { OE0E }
7278     { AF } { OE0F }
7279     { B0 } { OE10 }
7280     { B1 } { OE11 }
7281     { B2 } { OE12 }
7282     { B3 } { OE13 }
7283     { B4 } { OE14 }
7284     { B5 } { OE15 }
7285     { B6 } { OE16 }
7286     { B7 } { OE17 }
7287     { B8 } { OE18 }
7288     { B9 } { OE19 }
7289     { BA } { OE1A }
7290     { BB } { OE1B }
7291     { BC } { OE1C }
7292     { BD } { OE1D }
7293     { BE } { OE1E }
7294     { BF } { OE1F }
7295     { C0 } { OE20 }
7296     { C1 } { OE21 }
7297     { C2 } { OE22 }
7298     { C3 } { OE23 }
7299     { C4 } { OE24 }
7300     { C5 } { OE25 }
7301     { C6 } { OE26 }
7302     { C7 } { OE27 }
7303     { C8 } { OE28 }
7304     { C9 } { OE29 }
7305     { CA } { OE2A }
7306     { CB } { OE2B }
7307     { CC } { OE2C }
7308     { CD } { OE2D }
7309     { CE } { OE2E }
7310     { CF } { OE2F }
7311     { D0 } { OE30 }
7312     { D1 } { OE31 }
7313     { D2 } { OE32 }
7314     { D3 } { OE33 }
7315     { D4 } { OE34 }

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7316     { D5 } { 0E35 }
7317     { D6 } { 0E36 }
7318     { D7 } { 0E37 }
7319     { D8 } { 0E38 }
7320     { D9 } { 0E39 }
7321     { DA } { 0E3A }
7322     { DF } { 0E3F }
7323     { E0 } { 0E40 }
7324     { E1 } { 0E41 }
7325     { E2 } { 0E42 }
7326     { E3 } { 0E43 }
7327     { E4 } { 0E44 }
7328     { E5 } { 0E45 }
7329     { E6 } { 0E46 }
7330     { E7 } { 0E47 }
7331     { E8 } { 0E48 }
7332     { E9 } { 0E49 }
7333     { EA } { 0E4A }
7334     { EB } { 0E4B }
7335     { EC } { 0E4C }
7336     { ED } { 0E4D }
7337     { EE } { 0E4E }
7338     { EF } { 0E4F }
7339     { F0 } { 0E50 }
7340     { F1 } { 0E51 }
7341     { F2 } { 0E52 }
7342     { F3 } { 0E53 }
7343     { F4 } { 0E54 }
7344     { F5 } { 0E55 }
7345     { F6 } { 0E56 }
7346     { F7 } { 0E57 }
7347     { F8 } { 0E58 }
7348     { F9 } { 0E59 }
7349     { FA } { 0E5A }
7350     { FB } { 0E5B }
7351 }
7352 {
7353     { DB }
7354     { DC }
7355     { DD }
7356     { DE }
7357 }
7358 </iso885911>
7359 <iso885913>
7360 \__str_declare_eight_bit_encoding:nmmn { iso885913 } { 399 }
7361 {
7362     { A1 } { 201D }
7363     { A5 } { 201E }
7364     { A8 } { 00D8 }
7365     { AA } { 0156 }
7366     { AF } { 00C6 }
7367     { B4 } { 201C }
7368     { B8 } { 00F8 }
7369     { BA } { 0157 }

```

```

7370     { BF } { 00E6 }
7371     { C0 } { 0104 }
7372     { C1 } { 012E }
7373     { C2 } { 0100 }
7374     { C3 } { 0106 }
7375     { C6 } { 0118 }
7376     { C7 } { 0112 }
7377     { C8 } { 010C }
7378     { CA } { 0179 }
7379     { CB } { 0116 }
7380     { CC } { 0122 }
7381     { CD } { 0136 }
7382     { CE } { 012A }
7383     { CF } { 013B }
7384     { D0 } { 0160 }
7385     { D1 } { 0143 }
7386     { D2 } { 0145 }
7387     { D4 } { 014C }
7388     { D8 } { 0172 }
7389     { D9 } { 0141 }
7390     { DA } { 015A }
7391     { DB } { 016A }
7392     { DD } { 017B }
7393     { DE } { 017D }
7394     { E0 } { 0105 }
7395     { E1 } { 012F }
7396     { E2 } { 0101 }
7397     { E3 } { 0107 }
7398     { E6 } { 0119 }
7399     { E7 } { 0113 }
7400     { E8 } { 010D }
7401     { EA } { 017A }
7402     { EB } { 0117 }
7403     { EC } { 0123 }
7404     { ED } { 0137 }
7405     { EE } { 012B }
7406     { EF } { 013C }
7407     { F0 } { 0161 }
7408     { F1 } { 0144 }
7409     { F2 } { 0146 }
7410     { F4 } { 014D }
7411     { F8 } { 0173 }
7412     { F9 } { 0142 }
7413     { FA } { 015B }
7414     { FB } { 016B }
7415     { FD } { 017C }
7416     { FE } { 017E }
7417     { FF } { 2019 }
7418 }
7419 {
7420 }
7421 </iso885913>
7422 <*iso885914>
7423 \__str_declare_eight_bit_encoding:nnnn { iso885914 } { 529 }

```

```

7424 {
7425     { A1 } { 1E02 }
7426     { A2 } { 1E03 }
7427     { A4 } { 010A }
7428     { A5 } { 010B }
7429     { A6 } { 1E0A }
7430     { A8 } { 1E80 }
7431     { AA } { 1E82 }
7432     { AB } { 1E0B }
7433     { AC } { 1EF2 }
7434     { AF } { 0178 }
7435     { B0 } { 1E1E }
7436     { B1 } { 1E1F }
7437     { B2 } { 0120 }
7438     { B3 } { 0121 }
7439     { B4 } { 1E40 }
7440     { B5 } { 1E41 }
7441     { B7 } { 1E56 }
7442     { B8 } { 1E81 }
7443     { B9 } { 1E57 }
7444     { BA } { 1E83 }
7445     { BB } { 1E60 }
7446     { BC } { 1EF3 }
7447     { BD } { 1E84 }
7448     { BE } { 1E85 }
7449     { BF } { 1E61 }
7450     { D0 } { 0174 }
7451     { D7 } { 1E6A }
7452     { DE } { 0176 }
7453     { F0 } { 0175 }
7454     { F7 } { 1E6B }
7455     { FE } { 0177 }
7456 }
7457 {
7458 }
7459 </iso885914>
7460 <*iso885915>
7461 \_str_declare\_eight\_bit\_encoding:nnnn { iso885915 } { 383 }
7462 {
7463     { A4 } { 20AC }
7464     { A6 } { 0160 }
7465     { A8 } { 0161 }
7466     { B4 } { 017D }
7467     { B8 } { 017E }
7468     { BC } { 0152 }
7469     { BD } { 0153 }
7470     { BE } { 0178 }
7471 }
7472 {
7473 }
7474 </iso885915>
7475 <*iso885916>
7476 \_str_declare\_eight\_bit\_encoding:nnnn { iso885916 } { 558 }

```

```

7477 {
7478   { A1 } { 0104 }
7479   { A2 } { 0105 }
7480   { A3 } { 0141 }
7481   { A4 } { 20AC }
7482   { A5 } { 201E }
7483   { A6 } { 0160 }
7484   { A8 } { 0161 }
7485   { AA } { 0218 }
7486   { AC } { 0179 }
7487   { AE } { 017A }
7488   { AF } { 017B }
7489   { B2 } { 010C }
7490   { B3 } { 0142 }
7491   { B4 } { 017D }
7492   { B5 } { 201D }
7493   { B8 } { 017E }
7494   { B9 } { 010D }
7495   { BA } { 0219 }
7496   { BC } { 0152 }
7497   { BD } { 0153 }
7498   { BE } { 0178 }
7499   { BF } { 017C }
7500   { C3 } { 0102 }
7501   { C5 } { 0106 }
7502   { D0 } { 0110 }
7503   { D1 } { 0143 }
7504   { D5 } { 0150 }
7505   { D7 } { 015A }
7506   { D8 } { 0170 }
7507   { DD } { 0118 }
7508   { DE } { 021A }
7509   { E3 } { 0103 }
7510   { E5 } { 0107 }
7511   { F0 } { 0111 }
7512   { F1 } { 0144 }
7513   { F5 } { 0151 }
7514   { F7 } { 015B }
7515   { F8 } { 0171 }
7516   { FD } { 0119 }
7517   { FE } { 021B }
7518 }
7519 {
7520 }
7521 </iso885916>

```

11 l3seq implementation

The following test files are used for this code: m3seq002,m3seq003.

```

7522 <*package>
7523 <@@=seq>

```

A sequence is a control sequence whose top-level expansion is of the form “`\s__seq __seq_item:n {⟨item1⟩} ... __seq_item:n {⟨itemn⟩}`”, with a leading scan mark followed by n items of the same form. An earlier implementation used the structure “`\seq_elt:w ⟨item1⟩ \seq_elt_end: ... \seq_elt:w ⟨itemn⟩ \seq_elt_end:`”. This allowed rapid searching using a delimited function, but was not suitable for items containing `{`, `}` and `#` tokens, and also lead to the loss of surrounding braces around items

`__seq_item:n *`
`__seq_item:n {⟨item⟩}`

The internal token used to begin each sequence entry. If expanded outside of a mapping or manipulation function, an error is raised. The definition should always be set globally.

`__seq_push_item_def:n`
`__seq_push_item_def:x`
`__seq_push_item_def:n {⟨code⟩}`

Saves the definition of `__seq_item:n` and redefines it to accept one parameter and expand to `⟨code⟩`. This function should always be balanced by use of `__seq_pop_item_def:`.

`__seq_pop_item_def:`
`__seq_pop_item_def:`

Restores the definition of `__seq_item:n` most recently saved by `__seq_push_item_def:n`. This function should always be used in a balanced pair with `__seq_push_item_def:n`.

`\s__seq`

This private scan mark.

`7524 \scan_new:N \s__seq`

(End definition for `\s__seq`.)

`\s__seq_mark`

Private scan marks.

`\s__seq_stop`
`7525 \scan_new:N \s__seq_mark`
`7526 \scan_new:N \s__seq_stop`

(End definition for `\s__seq_mark` and `\s__seq_stop`.)

`__seq_item:n`

The delimiter is always defined, but when used incorrectly simply removes its argument and hits an undefined control sequence to raise an error.

`7527 \cs_new:Npn __seq_item:n`
`7528 {`
`7529 __kernel_msg_expandable_error:nn { kernel } { misused-sequence }`
`7530 \use_none:n`
`7531 }`

(End definition for `__seq_item:n`.)

`\l__seq_internal_a_tl`

Scratch space for various internal uses.

`\l__seq_internal_b_tl`
`7532 \tl_new:N \l__seq_internal_a_tl`
`7533 \tl_new:N \l__seq_internal_b_tl`

(End definition for `\l__seq_internal_a_tl` and `\l__seq_internal_b_tl`.)

`__seq_tmp:w`

Scratch function for internal use.

`7534 \cs_new_eq:NN __seq_tmp:w ?`

(End definition for `__seq_tmp:w`.)

\c_empty_seq A sequence with no item, following the structure mentioned above.

```
7535 \tl_const:Nn \c_empty_seq { \s__seq }
```

(End definition for \c_empty_seq. This variable is documented on page 86.)

11.1 Allocation and initialisation

\seq_new:N Sequences are initialized to \c_empty_seq.

```
\seq_new:c      7536 \cs_new_protected:Npn \seq_new:N #1
                  7537 {
                  7538   \__kernel_chk_if_free_cs:N #1
                  7539   \cs_gset_eq:NN #1 \c_empty_seq
                  7540 }
                  7541 \cs_generate_variant:Nn \seq_new:N { c }
```

(End definition for \seq_new:N. This function is documented on page 75.)

\seq_clear:N Clearing a sequence is similar to setting it equal to the empty one.

```
\seq_clear:c    7542 \cs_new_protected:Npn \seq_clear:N #1
\seq_gclear:N    7543 { \seq_set_eq:NN #1 \c_empty_seq }
\seq_gclear:c    7544 \cs_generate_variant:Nn \seq_clear:N { c }
                  7545 \cs_new_protected:Npn \seq_gclear:N #1
                  7546 { \seq_gset_eq:NN #1 \c_empty_seq }
                  7547 \cs_generate_variant:Nn \seq_gclear:N { c }
```

(End definition for \seq_clear:N and \seq_gclear:N. These functions are documented on page 75.)

\seq_clear_new:N Once again we copy code from the token list functions.

```
\seq_clear_new:c 7548 \cs_new_protected:Npn \seq_clear_new:N #1
\seq_gclear_new:N 7549 { \seq_if_exist:NTF #1 { \seq_clear:N #1 } { \seq_new:N #1 } }
\seq_gclear_new:c 7550 \cs_generate_variant:Nn \seq_clear_new:N { c }
                  7551 \cs_new_protected:Npn \seq_gclear_new:N #1
                  7552 { \seq_if_exist:NTF #1 { \seq_gclear:N #1 } { \seq_new:N #1 } }
                  7553 \cs_generate_variant:Nn \seq_gclear_new:N { c }
```

(End definition for \seq_clear_new:N and \seq_gclear_new:N. These functions are documented on page 75.)

\seq_set_eq:NN Copying a sequence is the same as copying the underlying token list.

```
\seq_set_eq:cN   7554 \cs_new_eq:NN \seq_set_eq:NN \tl_set_eq:NN
\seq_set_eq:Nc   7555 \cs_new_eq:NN \seq_set_eq:Nc \tl_set_eq:Nc
\seq_set_eq:cc   7556 \cs_new_eq:NN \seq_set_eq:cN \tl_set_eq:cN
\seq_gset_eq:NN  7557 \cs_new_eq:NN \seq_set_eq:cc \tl_set_eq:cc
\seq_gset_eq:cN  7558 \cs_new_eq:NN \seq_gset_eq:NN \tl_gset_eq:NN
\seq_gset_eq:Nc  7559 \cs_new_eq:NN \seq_gset_eq:Nc \tl_gset_eq:Nc
\seq_gset_eq:cN  7560 \cs_new_eq:NN \seq_gset_eq:cN \tl_gset_eq:cN
\seq_gset_eq:cc  7561 \cs_new_eq:NN \seq_gset_eq:cc \tl_gset_eq:cc
```

(End definition for \seq_set_eq:NN and \seq_gset_eq:NN. These functions are documented on page 75.)

`\seq_set_from_clist:NN` Setting a sequence from a comma-separated list is done using a simple mapping.

```

\seq_set_from_clist:cN 7562 \cs_new_protected:Npn \seq_set_from_clist:NN #1#2
\seq_set_from_clist:Nc 7563 {
\seq_set_from_clist:cc 7564   \__kernel_tl_set:Nx #1
\seq_set_from_clist:Nn 7565   { \s__seq \clist_map_function:NN #2 \__seq_wrap_item:n }
\seq_set_from_clist:cn 7566 }
\seq_gset_from_clist:NN 7567 \cs_new_protected:Npn \seq_gset_from_clist:Nn #1#2
\seq_gset_from_clist:cN 7568 {
\seq_gset_from_clist:Nc 7569   \__kernel_tl_set:Nx #1
\seq_gset_from_clist:cc 7570   { \s__seq \clist_map_function:nN {#2} \__seq_wrap_item:n }
\seq_gset_from_clist:Nn 7571 }
\seq_gset_from_clist:cn 7572 \cs_new_protected:Npn \seq_gset_from_clist:NN #1#2
7573 {
7574   \__kernel_tl_gset:Nx #1
7575   { \s__seq \clist_map_function:NN #2 \__seq_wrap_item:n }
7576 }
7577 \cs_new_protected:Npn \seq_gset_from_clist:Nn #1#2
7578 {
7579   \__kernel_tl_gset:Nx #1
7580   { \s__seq \clist_map_function:nN {#2} \__seq_wrap_item:n }
7581 }
7582 \cs_generate_variant:Nn \seq_set_from_clist:NN { Nc }
7583 \cs_generate_variant:Nn \seq_set_from_clist:NN { c , cc }
7584 \cs_generate_variant:Nn \seq_set_from_clist:Nn { c }
7585 \cs_generate_variant:Nn \seq_gset_from_clist:NN { Nc }
7586 \cs_generate_variant:Nn \seq_gset_from_clist:NN { c , cc }
7587 \cs_generate_variant:Nn \seq_gset_from_clist:Nn { c }

```

(End definition for `\seq_set_from_clist:NN` and others. These functions are documented on page 75.)

`\seq_const_from_clist:Nn` Almost identical to `\seq_set_from_clist:Nn`.

```

\seq_const_from_clist:cn 7588 \cs_new_protected:Npn \seq_const_from_clist:Nn #1#2
7589 {
7590   \tl_const:Nx #1
7591   { \s__seq \clist_map_function:nN {#2} \__seq_wrap_item:n }
7592 }
7593 \cs_generate_variant:Nn \seq_const_from_clist:Nn { c }

```

(End definition for `\seq_const_from_clist:Nn`. This function is documented on page 76.)

`\seq_set_split:Nnn` When the separator is empty, everything is very simple, just map `__seq_wrap_item:n` through the items of the last argument. For non-trivial separators, the goal is to split a given token list at the marker, strip spaces from each item, and remove one set of outer braces if after removing leading and trailing spaces the item is enclosed within braces. After `\tl_replace_all:Nnn`, the token list `\l__seq_internal_a_tl` is a repetition of the pattern `__seq_set_split_auxi:w \prg_do_nothing: <item with spaces> __seq_set_split_end:.` Then, x-expansion causes `__seq_set_split_auxi:w` to trim spaces, and leaves its result as `__seq_set_split_auxii:w <trimmed item> __seq_set_split_end:.` This is then converted to the `l3seq` internal structure by another x-expansion. In the first step, we insert `\prg_do_nothing:` to avoid losing braces too early; that would cause space trimming to act within those lost braces. The second step is solely there to strip braces which are outermost after space trimming.

```

7594 \cs_new_protected:Npn \seq_set_split:Nnn

```

```

7595 { \__seq_set_split:NNnn \__kernel_tl_set:Nx }
7596 \cs_new_protected:Npn \seq_gset_split:Nnn
7597 { \__seq_set_split:NNnn \__kernel_tl_gset:Nx }
7598 \cs_new_protected:Npn \__seq_set_split:NNnn #1#2#3#4
7599 {
7600   \tl_if_empty:nTF {#3}
7601   {
7602     \tl_set:Nn \l__seq_internal_a_tl
7603     { \tl_map_function:nN {#4} \__seq_wrap_item:n }
7604   }
7605   {
7606     \tl_set:Nn \l__seq_internal_a_tl
7607     {
7608       \__seq_set_split_auxi:w \prg_do_nothing:
7609       #4
7610       \__seq_set_split_end:
7611     }
7612     \tl_replace_all:Nnn \l__seq_internal_a_tl { #3 }
7613     {
7614       \__seq_set_split_end:
7615       \__seq_set_split_auxi:w \prg_do_nothing:
7616     }
7617     \__kernel_tl_set:Nx \l__seq_internal_a_tl { \l__seq_internal_a_tl }
7618   }
7619   #1 #2 { \s__seq \l__seq_internal_a_tl }
7620 }
7621 \cs_new:Npn \__seq_set_split_auxi:w #1 \__seq_set_split_end:
7622 {
7623   \exp_not:N \__seq_set_split_auxii:w
7624   \exp_args:No \tl_trim_spaces:n {#1}
7625   \exp_not:N \__seq_set_split_end:
7626 }
7627 \cs_new:Npn \__seq_set_split_auxii:w #1 \__seq_set_split_end:
7628 { \__seq_wrap_item:n {#1} }
7629 \cs_generate_variant:Nn \seq_set_split:Nnn { NnV }
7630 \cs_generate_variant:Nn \seq_gset_split:Nnn { NnV }

```

(End definition for \seq_set_split:Nnn and others. These functions are documented on page 76.)

\seq_concat:NNN When concatenating sequences, one must remove the leading \s__seq of the second sequence. The result starts with \s__seq (of the first sequence), which stops f-expansion.

\seq_concat:ccc

\seq_gconcat:NNN

```

7631 \cs_new_protected:Npn \seq_concat:NNN #1#2#3
7632 { \tl_set:Nf #1 { \exp_after:wN \use_i:nn \exp_after:wN #2 #3 } }
7633 \cs_new_protected:Npn \seq_gconcat:NNN #1#2#3
7634 { \tl_gset:Nf #1 { \exp_after:wN \use_i:nn \exp_after:wN #2 #3 } }
7635 \cs_generate_variant:Nn \seq_concat:NNN { ccc }
7636 \cs_generate_variant:Nn \seq_gconcat:NNN { ccc }

```

(End definition for \seq_concat:NNN and \seq_gconcat:NNN. These functions are documented on page 76.)

\seq_if_exist_p:N Copies of the cs functions defined in l3basics.

\seq_if_exist_p:c

```

7637 \prg_new_eq_conditional:NNn \seq_if_exist:N \cs_if_exist:N
7638 { TF , T , F , p }
7639 \prg_new_eq_conditional:NNn \seq_if_exist:c \cs_if_exist:c

```

\seq_if_exist:N \underline{TF}

\seq_if_exist:c \underline{TF}

```
7640 { TF , T , F , p }
```

(End definition for `\seq_if_exist:NTF`. This function is documented on page 76.)

11.2 Appending data to either end

`\seq_put_left:Nn` When adding to the left of a sequence, remove `\s__seq`. This is done by `__seq_put_left_aux:w`, which also stops f-expansion.

```

\seq_put_left:Nv 7641 \cs_new_protected:Npn \seq_put_left:Nn #1#2
\seq_put_left:No 7642 {
\seq_put_left:Nx 7643   \__kernel_tl_set:Nx #1
\seq_put_left:cn 7644   {
\seq_put_left:cV 7645     \exp_not:n { \s__seq \__seq_item:n {#2} }
\seq_put_left:cv 7646     \exp_not:f { \exp_after:wN \__seq_put_left_aux:w #1 }
\seq_put_left:co 7647   }
\seq_put_left:cx 7648 }
\seq_gput_left:Nn 7649 \cs_new_protected:Npn \seq_gput_left:Nn #1#2
\seq_gput_left:Nv 7650 {
\seq_gput_left:Nx 7651   \__kernel_tl_gset:Nx #1
\seq_gput_left:No 7652   {
\seq_gput_left:Nx 7653     \exp_not:n { \s__seq \__seq_item:n {#2} }
\seq_gput_left:cn 7654     \exp_not:f { \exp_after:wN \__seq_put_left_aux:w #1 }
\seq_gput_left:cV 7655   }
\seq_gput_left:cv 7656 }
\seq_gput_left:co 7657 \cs_new:Npn \__seq_put_left_aux:w \s__seq { \exp_stop_f: }
\seq_gput_left:cx 7658 \cs_generate_variant:Nn \seq_put_left:Nn { NV , Nv , No , Nx }
\__seq_put_left_aux:w 7659 \cs_generate_variant:Nn \seq_put_left:Nn { c , cV , cv , co , cx }
7660 \cs_generate_variant:Nn \seq_gput_left:Nn { NV , Nv , No , Nx }
7661 \cs_generate_variant:Nn \seq_gput_left:Nn { c , cV , cv , co , cx }
```

(End definition for `\seq_put_left:Nn`, `\seq_gput_left:Nn`, and `__seq_put_left_aux:w`. These functions are documented on page 76.)

`\seq_put_right:Nn` Since there is no trailing marker, adding an item to the right of a sequence simply means wrapping it in `__seq_item:n`.

```

\seq_put_right:Nv 7662 \cs_new_protected:Npn \seq_put_right:Nn #1#2
\seq_put_right:No 7663 { \tl_put_right:Nn #1 { \__seq_item:n {#2} } }
\seq_put_right:Nx 7664 \cs_new_protected:Npn \seq_gput_right:Nn #1#2
\seq_put_right:cn 7665 { \tl_gput_right:Nn #1 { \__seq_item:n {#2} } }
\seq_put_right:cV 7666 \cs_generate_variant:Nn \seq_gput_right:Nn { NV , Nv , No , Nx }
\seq_put_right:cv 7667 \cs_generate_variant:Nn \seq_gput_right:Nn { c , cV , cv , co , cx }
\seq_put_right:co 7668 \cs_generate_variant:Nn \seq_put_right:Nn { NV , Nv , No , Nx }
\seq_put_right:cx 7669 \cs_generate_variant:Nn \seq_put_right:Nn { c , cV , cv , co , cx }
```

`\seq_gput_right:Nn` (End definition for `\seq_put_right:Nn` and `\seq_gput_right:Nn`. These functions are documented on page 76.)

```

\seq_gput_right:Nv
\seq_gput_right:No
\seq_gput_right:Nx
\seq_gput_right:cn
\seq_gput_right:cV
\seq_gput_right:cv
\seq_gput_right:co
\seq_gput_right:cx
```

11.3 Modifying sequences

This function converts its argument to a proper sequence item in an x-expansion context.

```
7670 \cs_new:Npn \__seq_wrap_item:n #1 { \exp_not:n { \__seq_item:n {#1} } }
```

(End definition for `__seq_wrap_item:n`.)

`\l__seq_remove_seq` An internal sequence for the removal routines.

```
7671 \seq_new:N \l__seq_remove_seq
```

(End definition for `\l__seq_remove_seq`.)

`\seq_remove_duplicates:N` Removing duplicates means making a new list then copying it.

```
\seq_remove_duplicates:c 7672 \cs_new_protected:Npn \seq_remove_duplicates:N
\seq_gremove_duplicates:N 7673 { \__seq_remove_duplicates:NN \seq_set_eq:NN }
\seq_gremove_duplicates:c 7674 \cs_new_protected:Npn \seq_gremove_duplicates:N
\__seq_remove_duplicates:NN 7675 { \__seq_remove_duplicates:NN \seq_gset_eq:NN }
7676 \cs_new_protected:Npn \__seq_remove_duplicates:NN #1#2
7677 {
7678   \seq_clear:N \l__seq_remove_seq
7679   \seq_map_inline:Nn #2
7680   {
7681     \seq_if_in:NnF \l__seq_remove_seq {##1}
7682     { \seq_put_right:Nn \l__seq_remove_seq {##1} }
7683   }
7684   #1 #2 \l__seq_remove_seq
7685 }
7686 \cs_generate_variant:Nn \seq_remove_duplicates:N { c }
7687 \cs_generate_variant:Nn \seq_gremove_duplicates:N { c }
```

(End definition for `\seq_remove_duplicates:N`, `\seq_gremove_duplicates:N`, and `__seq_remove_duplicates:NN`. These functions are documented on page 79.)

`\seq_remove_all:Nn` The idea of the code here is to avoid a relatively expensive addition of items one at a time
`\seq_remove_all:cn` to an intermediate sequence. The approach taken is therefore similar to that in `__seq_`
`\seq_gremove_all:Nn` `pop_right:NNN`, using a “flexible” x-type expansion to do most of the work. As `\tl_`
`\seq_gremove_all:cn` `if_eq:nnT` is not expandable, a two-part strategy is needed. First, the x-type expansion
`__seq_remove_all_aux:NNn` uses `\str_if_eq:nnT` to find potential matches. If one is found, the expansion is halted
and the necessary set up takes place to use the `\tl_if_eq:NNT` test. The x-type is started
again, including all of the items copied already. This happens repeatedly until the entire
sequence has been scanned. The code is set up to avoid needing and intermediate scratch
list: the lead-off x-type expansion (`#1 #2 {#2}`) ensures that nothing is lost.

```
7688 \cs_new_protected:Npn \seq_remove_all:Nn
7689 { \__seq_remove_all_aux:NNn \__kernel_tl_set:Nx }
7690 \cs_new_protected:Npn \seq_gremove_all:Nn
7691 { \__seq_remove_all_aux:NNn \__kernel_tl_gset:Nx }
7692 \cs_new_protected:Npn \__seq_remove_all_aux:NNn #1#2#3
7693 {
7694   \__seq_push_item_def:n
7695   {
7696     \str_if_eq:nnT {##1} {#3}
7697     {
7698       \if_false: { \fi: }
7699       \tl_set:Nn \l__seq_internal_b_tl {##1}
7700       #1 #2
7701       { \if_false: } \fi:
7702       \exp_not:o {#2}
7703       \tl_if_eq:NNT \l__seq_internal_a_tl \l__seq_internal_b_tl
7704       { \use_none:nn }
7705     }
7706     \__seq_wrap_item:n {##1}
```

```

7707     }
7708     \tl_set:Nn \l__seq_internal_a_tl {#3}
7709     #1 #2 {#2}
7710     \__seq_pop_item_def:
7711   }
7712   \cs_generate_variant:Nn \seq_remove_all:Nn { c }
7713   \cs_generate_variant:Nn \seq_gremove_all:Nn { c }

```

(End definition for `\seq_remove_all:Nn`, `\seq_gremove_all:Nn`, and `__seq_remove_all_aux:NNn`. These functions are documented on page 79.)

```

\seq_reverse:N Previously, \seq_reverse:N was coded by collecting the items in reverse order after an
\seq_reverse:c \exp_stop_f: marker.
\seq_greverse:N \cs_new_protected:Npn \seq_reverse:N #1
\seq_greverse:c {
  \cs_set_eq:NN \@@_item:n \@@_reverse_item:nw
  \tl_set:Nf #2 { #2 \exp_stop_f: }
}
\cs_new:Npn \@@_reverse_item:nw #1 #2 \exp_stop_f:
{
  #2 \exp_stop_f:
  \@@_item:n {#1}
}

```

At first, this seems optimal, since we can forget about each item as soon as it is placed after `\exp_stop_f:`. Unfortunately, \TeX 's usual tail recursion does not take place in this case: since the following `__seq_reverse_item:nw` only reads tokens until `\exp_stop_f:`, and never reads the `\@@_item:n {#1}` left by the previous call, \TeX cannot remove that previous call from the stack, and in particular must retain the various macro parameters in memory, until the end of the replacement text is reached. The stack is thus only flushed after all the `__seq_reverse_item:nw` are expanded. Keeping track of the arguments of all those calls uses up a memory quadratic in the length of the sequence. \TeX can then not cope with more than a few thousand items.

Instead, we collect the items in the argument of `\exp_not:n`. The previous calls are cleanly removed from the stack, and the memory consumption becomes linear.

```

7714 \cs_new_protected:Npn \seq_reverse:N
7715 { \__seq_reverse:NN \__kernel_tl_set:Nx }
7716 \cs_new_protected:Npn \seq_greverse:N
7717 { \__seq_reverse:NN \__kernel_tl_gset:Nx }
7718 \cs_new_protected:Npn \__seq_reverse:NN #1 #2
7719 {
7720   \cs_set_eq:NN \__seq_tmp:w \__seq_item:n
7721   \cs_set_eq:NN \__seq_item:n \__seq_reverse_item:nwn
7722   #1 #2 { #2 \exp_not:n { } }
7723   \cs_set_eq:NN \__seq_item:n \__seq_tmp:w
7724 }
7725 \cs_new:Npn \__seq_reverse_item:nwn #1 #2 \exp_not:n #3
7726 {
7727   #2
7728   \exp_not:n { \__seq_item:n {#1} #3 }
7729 }
7730 \cs_generate_variant:Nn \seq_reverse:N { c }
7731 \cs_generate_variant:Nn \seq_greverse:N { c }

```

(End definition for `\seq_reverse:N` and others. These functions are documented on page 79.)

`\seq_sort:Nn` Implemented in `l3sort`.

`\seq_sort:cn`

`\seq_gsort:Nn` (End definition for `\seq_sort:Nn` and `\seq_gsort:Nn`. These functions are documented on page 79.)

`\seq_gsort:cn`

11.4 Sequence conditionals

`\seq_if_empty_p:N` Similar to token lists, we compare with the empty sequence.

```

7732 \prg_new_conditional:Npnn \seq_if_empty:N #1 { p , T , F , TF }
7733 {
7734   \if_meaning:w #1 \c_empty_seq
7735   \prg_return_true:
7736   \else:
7737     \prg_return_false:
7738   \fi:
7739 }
7740 \prg_generate_conditional_variant:Nnn \seq_if_empty:N
7741 { c } { p , T , F , TF }
```

(End definition for `\seq_if_empty:N`. This function is documented on page 80.)

`\seq_shuffle:N` We apply the Fisher–Yates shuffle, storing items in `\toks` registers. We use the primitive
`\seq_shuffle:c` `\tex_uniformdeviate:D` for speed reasons. Its non-uniformity is of order its argument
`\seq_gshuffle:N` divided by 2^{28} , not too bad for small lists. For sequences with more than 13 elements
`\seq_gshuffle:c` there are more possible permutations than possible seeds ($13! > 2^{28}$) so the question
`__seq_shuffle:NN` of uniformity is somewhat moot. The integer variables are declared in `l3int`: load-order
`__seq_shuffle_item:n` issues.

```

7742 \cs_if_exist:NTF \tex_uniformdeviate:D
7743 {
7744   \seq_new:N \g__seq_internal_seq
7745   \cs_new_protected:Npn \seq_shuffle:N { \__seq_shuffle:NN \seq_set_eq:NN }
7746   \cs_new_protected:Npn \seq_gshuffle:N { \__seq_shuffle:NN \seq_gset_eq:NN }
7747   \cs_new_protected:Npn \__seq_shuffle:NN #1#2
7748   {
7749     \int_compare:nNnTF { \seq_count:N #2 } > \c_max_register_int
7750     {
7751       \__kernel_msg_error:nxx { kernel } { shuffle-too-large }
7752       { \token_to_str:N #2 }
7753     }
7754     {
7755       \group_begin:
7756         \int_zero:N \l__seq_internal_a_int
7757         \__seq_push_item_def:
7758         \cs_gset_eq:NN \__seq_item:n \__seq_shuffle_item:n
7759         #2
7760         \__seq_pop_item_def:
7761         \seq_gset_from_inline_x:Nnn \g__seq_internal_seq
7762         { \int_step_function:nN { \l__seq_internal_a_int } }
7763         { \tex_the:D \tex_toks:D ##1 }
7764       \group_end:
7765       #1 #2 \g__seq_internal_seq
7766       \seq_gclear:N \g__seq_internal_seq
7767     }
7768   }
```

```

7768     }
7769 \cs_new_protected:Npn \__seq_shuffle_item:n
7770 {
7771     \int_incr:N \l__seq_internal_a_int
7772     \int_set:Nn \l__seq_internal_b_int
7773     { 1 + \tex_uniformdeviate:D \l__seq_internal_a_int }
7774     \tex_toks:D \l__seq_internal_a_int
7775     = \tex_toks:D \l__seq_internal_b_int
7776     \tex_toks:D \l__seq_internal_b_int
7777 }
7778 }
7779 {
7780 \cs_new_protected:Npn \seq_shuffle:N #1
7781 {
7782     \__kernel_msg_error:nnn { kernel } { fp-no-random }
7783     { \seq_shuffle:N #1 }
7784 }
7785 \cs_new_eq:NN \seq_gshuffle:N \seq_shuffle:N
7786 }
7787 \cs_generate_variant:Nn \seq_shuffle:N { c }
7788 \cs_generate_variant:Nn \seq_gshuffle:N { c }

```

(End definition for `\seq_shuffle:N` and others. These functions are documented on page 80.)

`\seq_if_in:NnTF` The approach here is to define `__seq_item:n` to compare its argument with the test sequence. If the two items are equal, the mapping is terminated and `\group_end: \prg_return_true:` is inserted after skipping over the rest of the recursion. On the other hand, if there is no match then the loop breaks, returning `\prg_return_false:`. Everything is inside a group so that `__seq_item:n` is preserved in nested situations.

```

\seq_if_in:NvTF
\seq_if_in:NoTF
\seq_if_in:NxTF
\seq_if_in:cnTF
\seq_if_in:cVTF
\seq_if_in:cvTF
\seq_if_in:coTF
\seq_if_in:cxTF
\__seq_if_in:
7789 \prg_new_protected_conditional:Npnn \seq_if_in:Nn #1#2
7790 { T , F , TF }
7791 {
7792     \group_begin:
7793     \tl_set:Nn \l__seq_internal_a_tl {#2}
7794     \cs_set_protected:Npn \__seq_item:n ##1
7795     {
7796         \tl_set:Nn \l__seq_internal_b_tl {##1}
7797         \if_meaning:w \l__seq_internal_a_tl \l__seq_internal_b_tl
7798         \exp_after:wN \__seq_if_in:
7799         \fi:
7800     }
7801     #1
7802     \group_end:
7803     \prg_return_false:
7804     \prg_break_point:
7805 }
7806 \cs_new:Npn \__seq_if_in:
7807 { \prg_break:n { \group_end: \prg_return_true: } }
7808 \prg_generate_conditional_variant:Nnn \seq_if_in:Nn
7809 { NV , Nv , No , Nx , c , cV , cv , co , cx } { T , F , TF }

```

(End definition for `\seq_if_in:NnTF` and `__seq_if_in:`. This function is documented on page 80.)

11.5 Recovering data from sequences

`__seq_pop:NNNN` The two `pop` functions share their emptiness tests. We also use a common emptiness test
`__seq_pop_TF:NNNN` for all branching `get` and `pop` functions.

```

7810 \cs_new_protected:Npn \__seq_pop:NNNN #1#2#3#4
7811 {
7812   \if_meaning:w #3 \c_empty_seq
7813     \tl_set:Nn #4 { \q_no_value }
7814   \else:
7815     #1#2#3#4
7816   \fi:
7817 }
7818 \cs_new_protected:Npn \__seq_pop_TF:NNNN #1#2#3#4
7819 {
7820   \if_meaning:w #3 \c_empty_seq
7821     % \tl_set:Nn #4 { \q_no_value }
7822     \prg_return_false:
7823   \else:
7824     #1#2#3#4
7825     \prg_return_true:
7826   \fi:
7827 }
```

(End definition for `__seq_pop:NNNN` and `__seq_pop_TF:NNNN`.)

`\seq_get_left:NN` Getting an item from the left of a sequence is pretty easy: just trim off the first item
`\seq_get_left:cN` after `__seq_item:n` at the start. We append a `\q_no_value` item to cover the case of
`__seq_get_left:wnw` an empty sequence

```

7828 \cs_new_protected:Npn \seq_get_left:NN #1#2
7829 {
7830   \__kernel_tl_set:Nx #2
7831   {
7832     \exp_after:wN \__seq_get_left:wnw
7833     #1 \__seq_item:n { \q_no_value } \s__seq_stop
7834   }
7835 }
7836 \cs_new:Npn \__seq_get_left:wnw #1 \__seq_item:n #2#3 \s__seq_stop
7837 { \exp_not:n {#2} }
7838 \cs_generate_variant:Nn \seq_get_left:NN { c }
```

(End definition for `\seq_get_left:NN` and `__seq_get_left:wnw`. This function is documented on page 77.)

`\seq_pop_left:NN` The approach to popping an item is pretty similar to that to get an item, with the only
`\seq_pop_left:cN` difference being that the sequence itself has to be redefined. This makes it more sensible
`\seq_gpop_left:NN` to use an auxiliary function for the local and global cases.
`\seq_gpop_left:cN`

```

7839 \cs_new_protected:Npn \seq_pop_left:NN
7840 { \__seq_pop:NNNN \__seq_pop_left:NNN \tl_set:Nn }
7841 \cs_new_protected:Npn \seq_gpop_left:NN
7842 { \__seq_pop:NNNN \__seq_pop_left:NNN \tl_gset:Nn }
7843 \cs_new_protected:Npn \__seq_pop_left:NNN #1#2#3
7844 { \exp_after:wN \__seq_pop_left:wnwNNN #2 \s__seq_stop #1#2#3 }
7845 \cs_new_protected:Npn \__seq_pop_left:wnwNNN
7846 #1 \__seq_item:n #2#3 \s__seq_stop #4#5#6
```

```

7847 {
7848     #4 #5 { #1 #3 }
7849     \tl_set:Nn #6 {#2}
7850 }
7851 \cs_generate_variant:Nn \seq_pop_left:NN { c }
7852 \cs_generate_variant:Nn \seq_gpop_left:NN { c }

```

(End definition for `\seq_pop_left:NN` and others. These functions are documented on page 77.)

`\seq_get_right:NN` First remove `\s__seq` and prepend `\q_no_value`. The first argument of `__seq_get_right_loop:nw` is the last item found, and the second argument is empty until the end of the loop, where it is code that applies `\exp_not:n` to the last item and ends the loop.

```

\__seq_get_right_loop:nw
\__seq_get_right_end:NnN
7853 \cs_new_protected:Npn \seq_get_right:NN #1#2
7854 {
7855     \__kernel_tl_set:Nx #2
7856     {
7857         \exp_after:wN \use_i_ii:nnn
7858         \exp_after:wN \__seq_get_right_loop:nw
7859         \exp_after:wN \q_no_value
7860         #1
7861         \__seq_get_right_end:NnN \__seq_item:n
7862     }
7863 }
7864 \cs_new:Npn \__seq_get_right_loop:nw #1#2 \__seq_item:n
7865 {
7866     #2 \use_none:n {#1}
7867     \__seq_get_right_loop:nw
7868 }
7869 \cs_new:Npn \__seq_get_right_end:NnN #1#2#3 { \exp_not:n {#2} }
7870 \cs_generate_variant:Nn \seq_get_right:NN { c }

```

(End definition for `\seq_get_right:NN`, `__seq_get_right_loop:nw`, and `__seq_get_right_end:NnN`. This function is documented on page 77.)

`\seq_pop_right:NN` The approach to popping from the right is a bit more involved, but does use some of the same ideas as getting from the right. What is needed is a “flexible length” way to set a token list variable. This is supplied by the `{\if_false:} \fi: \dots \if_false: { \fi: }` construct. Using an x-type expansion and a “non-expanding” definition for `__seq_item:n`, the left-most $n - 1$ entries in a sequence of n items are stored back in the sequence. That needs a loop of unknown length, hence using the strange `\if_false:` way of including braces. When the last item of the sequence is reached, the closing brace for the assignment is inserted, and `\tl_set:Nn #3` is inserted in front of the final entry. This therefore does the pop assignment. One more iteration is performed, with an empty argument and `\use_none:nn`, which finally stops the loop.

```

7871 \cs_new_protected:Npn \seq_pop_right:NN
7872 { \__seq_pop:NNNN \__seq_pop_right:NNN \__kernel_tl_set:Nx }
7873 \cs_new_protected:Npn \seq_gpop_right:NN
7874 { \__seq_pop:NNNN \__seq_pop_right:NNN \__kernel_tl_gset:Nx }
7875 \cs_new_protected:Npn \__seq_pop_right:NNN #1#2#3
7876 {
7877     \cs_set_eq:NN \__seq_tmp:w \__seq_item:n
7878     \cs_set_eq:NN \__seq_item:n \scan_stop:
7879     #1 #2
7880     { \if_false: } \fi: \s__seq

```

```

7881         \exp_after:wN \use_i:nnn
7882         \exp_after:wN \__seq_pop_right_loop:nn
7883         #2
7884         {
7885             \if_false: { \fi: }
7886             \__kernel_tl_set:Nx #3
7887         }
7888         { } \use_none:nn
7889         \cs_set_eq:NN \__seq_item:n \__seq_tmp:w
7890     }
7891 \cs_new:Npn \__seq_pop_right_loop:nn #1#2
7892 {
7893     #2 { \exp_not:n {#1} }
7894     \__seq_pop_right_loop:nn
7895 }
7896 \cs_generate_variant:Nn \seq_pop_right:NN { c }
7897 \cs_generate_variant:Nn \seq_gpop_right:NN { c }

```

(End definition for \seq_pop_right:NN and others. These functions are documented on page 77.)

\seq_get_left:NNTF Getting from the left or right with a check on the results. The first argument to __seq_pop_TF:NNNN is left unused.

```

7898 \seq_get_left:cNTF \prg_new_protected_conditional:Npnn \seq_get_left:NN #1#2 { T , F , TF }
7899 { \__seq_pop_TF:NNNN \prg_do_nothing: \seq_get_left:NN #1#2 }
7900 \seq_get_right:NNTF \prg_new_protected_conditional:Npnn \seq_get_right:NN #1#2 { T , F , TF }
7901 { \__seq_pop_TF:NNNN \prg_do_nothing: \seq_get_right:NN #1#2 }
7902 \prg_generate_conditional_variant:Nnn \seq_get_left:NN
7903 { c } { T , F , TF }
7904 \prg_generate_conditional_variant:Nnn \seq_get_right:NN
7905 { c } { T , F , TF }

```

(End definition for \seq_get_left:NNTF and \seq_get_right:NNTF. These functions are documented on page 78.)

\seq_pop_left:NNTF More or less the same for popping.

```

7906 \seq_pop_left:cNTF \prg_new_protected_conditional:Npnn \seq_pop_left:NN #1#2
7907 { T , F , TF }
7908 \seq_gpop_left:NNTF \prg_new_protected_conditional:Npnn \seq_gpop_left:NN #1#2
7909 { T , F , TF }
7910 \seq_pop_right:NNTF \prg_new_protected_conditional:Npnn \seq_gpop_left:NN #1#2
7911 { T , F , TF }
7912 \seq_gpop_right:NNTF \prg_new_protected_conditional:Npnn \seq_pop_right:NN #1#2
7913 { T , F , TF }
7914 \prg_generate_conditional_variant:Nnn \seq_pop_right:NN \__kernel_tl_set:Nx #1 #2 }
7915 \prg_generate_conditional_variant:Nnn \seq_gpop_right:NN #1#2
7916 { T , F , TF }
7917 \prg_generate_conditional_variant:Nnn \seq_pop_right:NN \__kernel_tl_gset:Nx #1 #2 }
7918 \prg_generate_conditional_variant:Nnn \seq_pop_left:NN { c }
7919 { T , F , TF }
7920 \prg_generate_conditional_variant:Nnn \seq_gpop_left:NN { c }
7921 { T , F , TF }
7922 \prg_generate_conditional_variant:Nnn \seq_pop_right:NN { c }
7923 { T , F , TF }
7924 \prg_generate_conditional_variant:Nnn \seq_gpop_right:NN { c }
7925 { T , F , TF }

```

(End definition for `\seq_pop_left:NNTF` and others. These functions are documented on page 78.)

`\seq_item:Nn` The idea here is to find the offset of the item from the left, then use a loop to grab the correct item. If the resulting offset is too large, then the argument delimited by `__seq_item:wNn` `\seq_item:n` is `\prg_break:` instead of being empty, terminating the loop and returning nothing at all.

```

7926 \cs_new:Npn \seq_item:Nn #1
7927   { \exp_after:wN \__seq_item:wNn #1 \s__seq_stop #1 }
7928 \cs_new:Npn \__seq_item:wNn \s__seq #1 \s__seq_stop #2#3
7929   {
7930     \exp_args:Nf \__seq_item:nwn
7931     { \exp_args:Nf \__seq_item:nN { \int_eval:n {#3} } #2 }
7932     #1
7933     \prg_break: \__seq_item:n { }
7934     \prg_break_point:
7935   }
7936 \cs_new:Npn \__seq_item:nN #1#2
7937   {
7938     \int_compare:nNnTF {#1} < 0
7939     { \int_eval:n { \seq_count:N #2 + 1 + #1 } }
7940     {#1}
7941   }
7942 \cs_new:Npn \__seq_item:nwn #1#2 \__seq_item:n #3
7943   {
7944     #2
7945     \int_compare:nNnTF {#1} = 1
7946     { \prg_break:n { \exp_not:n {#3} } }
7947     { \exp_args:Nf \__seq_item:nwn { \int_eval:n { #1 - 1 } } }
7948   }
7949 \cs_generate_variant:Nn \seq_item:Nn { c }

```

(End definition for `\seq_item:Nn` and others. This function is documented on page 77.)

`\seq_rand_item:N` Importantly, `\seq_item:Nn` only evaluates its argument once.

```

7950 \cs_new:Npn \seq_rand_item:N #1
7951   {
7952     \seq_if_empty:NF #1
7953     { \seq_item:Nn #1 { \int_rand:nn { 1 } { \seq_count:N #1 } } }
7954   }
7955 \cs_generate_variant:Nn \seq_rand_item:N { c }

```

(End definition for `\seq_rand_item:N`. This function is documented on page 78.)

11.6 Mapping to sequences

`\seq_map_break:` To break a function, the special token `\prg_break_point:Nn` is used to find the end of the code. Any ending code is then inserted before the return value of `\seq_map_break:n` is inserted.

```

7956 \cs_new:Npn \seq_map_break:
7957   { \prg_map_break:Nn \seq_map_break: { } }
7958 \cs_new:Npn \seq_map_break:n
7959   { \prg_map_break:Nn \seq_map_break: }

```

(End definition for `\seq_map_break:` and `\seq_map_break:n`. These functions are documented on page 82.)

`\seq_map_function:NN` The idea here is to apply the code of #2 to each item in the sequence without altering
`\seq_map_function:cN` the definition of `__seq_item:n`. The argument delimited by `__seq_item:n` is almost
`__seq_map_function:NNn` always empty, except at the end of the loop where it is `\prg_break:`. This allows to
break the loop without needing to do a (relatively-expensive) quark test.

```

7960 \cs_new:Npn \seq_map_function:NN #1#2
7961 {
7962   \exp_after:wN \use_i_ii:nnn
7963   \exp_after:wN \__seq_map_function:Nw
7964   \exp_after:wN #2
7965   #1
7966   \prg_break: \__seq_item:n { } \prg_break_point:
7967   \prg_break_point:Nn \seq_map_break: { }
7968 }
7969 \cs_new:Npn \__seq_map_function:Nw #1#2 \__seq_item:n #3
7970 {
7971   #2
7972   #1 {#3}
7973   \__seq_map_function:Nw #1
7974 }
7975 \cs_generate_variant:Nn \seq_map_function:NN { c }

```

(End definition for `\seq_map_function:NN` and `__seq_map_function:NNn`. This function is documented on page 80.)

`__seq_push_item_def:n` The definition of `__seq_item:n` needs to be saved and restored at various points within
`__seq_push_item_def:x` the mapping and manipulation code. That is handled here: as always, this approach uses
`__seq_push_item_def:` global assignments.
`__seq_pop_item_def:`

```

7976 \cs_new_protected:Npn \__seq_push_item_def:n
7977 {
7978   \__seq_push_item_def:
7979   \cs_gset:Npn \__seq_item:n ##1
7980 }
7981 \cs_new_protected:Npn \__seq_push_item_def:x
7982 {
7983   \__seq_push_item_def:
7984   \cs_gset:Npx \__seq_item:n ##1
7985 }
7986 \cs_new_protected:Npn \__seq_push_item_def:
7987 {
7988   \int_gincr:N \g__kernel_prg_map_int
7989   \cs_gset_eq:cN { __seq_map_ \int_use:N \g__kernel_prg_map_int :w }
7990   \__seq_item:n
7991 }
7992 \cs_new_protected:Npn \__seq_pop_item_def:
7993 {
7994   \cs_gset_eq:Nc \__seq_item:n
7995   { __seq_map_ \int_use:N \g__kernel_prg_map_int :w }
7996   \int_gdecr:N \g__kernel_prg_map_int
7997 }

```

(End definition for `__seq_push_item_def:n`, `__seq_push_item_def:`, and `__seq_pop_item_def:`.)

\seq_map_inline:Nn The idea here is that `__seq_item:n` is already “applied” to each item in a sequence, and so an in-line mapping is just a case of redefining `__seq_item:n`.

```

7998 \cs_new_protected:Npn \seq_map_inline:Nn #1#2
7999 {
8000     \__seq_push_item_def:n {#2}
8001     #1
8002     \prg_break_point:Nn \seq_map_break: { \__seq_pop_item_def: }
8003 }
8004 \cs_generate_variant:Nn \seq_map_inline:Nn { c }

```

(End definition for `\seq_map_inline:Nn`. This function is documented on page 80.)

\seq_map_tokens:Nn This is based on the function mapping but using the same tricks as described for `\prop_map_tokens:Nn`. The idea is to remove the leading `\s__seq` and apply the tokens such that they are safe with the break points, hence the `\use:n`.

```

8005 \cs_new:Npn \seq_map_tokens:Nn #1#2
8006 {
8007     \exp_last_unbraced:Nno
8008     \use_i:nn { \__seq_map_tokens:nw {#2} } #1
8009     \prg_break: \__seq_item:n { } \prg_break_point:
8010     \prg_break_point:Nn \seq_map_break: { }
8011 }
8012 \cs_generate_variant:Nn \seq_map_tokens:Nn { c }
8013 \cs_new:Npn \__seq_map_tokens:nw #1#2 \__seq_item:n #3
8014 {
8015     #2
8016     \use:n {#1} {#3}
8017     \__seq_map_tokens:nw {#1}
8018 }

```

(End definition for `\seq_map_tokens:Nn` and `__seq_map_tokens:nw`. This function is documented on page 81.)

\seq_map_variable:NNn This is just a specialised version of the in-line mapping function, using an `x`-type expansion for the code set up so that the number of `#` tokens required is as expected.

```

8019 \cs_new_protected:Npn \seq_map_variable:NNn #1#2#3
8020 {
8021     \__seq_push_item_def:x
8022     {
8023         \tl_set:Nn \exp_not:N #2 {##1}
8024         \exp_not:n {#3}
8025     }
8026     #1
8027     \prg_break_point:Nn \seq_map_break: { \__seq_pop_item_def: }
8028 }
8029 \cs_generate_variant:Nn \seq_map_variable:NNn { Nc }
8030 \cs_generate_variant:Nn \seq_map_variable:NNn { c , cc }

```

(End definition for `\seq_map_variable:NNn`. This function is documented on page 81.)

\seq_map_indexed_function:NN Similar to `\seq_map_function:NN` but we keep track of the item index as a `;`-delimited argument of `__seq_map_indexed:Nw`.

```

\seq_map_indexed_inline:Nn
\__seq_map_indexed:nNn
\__seq_map_indexed:Nw
8031 \cs_new:Npn \seq_map_indexed_function:NN #1#2
8032 {

```

```

8033     \__seq_map_indexed:NN #1#2
8034     \prg_break_point:Nn \seq_map_break: { }
8035 }
8036 \cs_new_protected:Npn \seq_map_indexed_inline:Nn #1#2
8037 {
8038     \int_gincr:N \g__kernel_prg_map_int
8039     \cs_gset_protected:cpn
8040     { \__seq_map_ \int_use:N \g__kernel_prg_map_int :w } ##1##2 {#2}
8041     \exp_args:NNc \__seq_map_indexed:NN #1
8042     { \__seq_map_ \int_use:N \g__kernel_prg_map_int :w }
8043     \prg_break_point:Nn \seq_map_break:
8044     { \int_gdecr:N \g__kernel_prg_map_int }
8045 }
8046 \cs_new:Npn \__seq_map_indexed:NN #1#2
8047 {
8048     \exp_after:wN \__seq_map_indexed:Nw
8049     \exp_after:wN #2
8050     \int_value:w 1
8051     \exp_after:wN \use_i:nn
8052     \exp_after:wN ;
8053     #1
8054     \prg_break: \__seq_item:n { } \prg_break_point:
8055 }
8056 \cs_new:Npn \__seq_map_indexed:Nw #1#2 ; #3 \__seq_item:n #4
8057 {
8058     #3
8059     #1 {#2} {#4}
8060     \exp_after:wN \__seq_map_indexed:Nw
8061     \exp_after:wN #1
8062     \int_value:w \int_eval:w 1 + #2 ;
8063 }

```

(End definition for `\seq_map_indexed_function:NN` and others. These functions are documented on page 81.)

`\seq_set_map_x:NNn`
`\seq_gset_map_x:NNn`
`__seq_set_map_x:NNNn`

Very similar to `\seq_set_filter:NNn`. We could actually merge the two within a single function, but it would have weird semantics.

```

8064 \cs_new_protected:Npn \seq_set_map_x:NNn
8065 { \__seq_set_map_x:NNNn \__kernel_tl_set:Nx }
8066 \cs_new_protected:Npn \seq_gset_map_x:NNn
8067 { \__seq_set_map_x:NNNn \__kernel_tl_gset:Nx }
8068 \cs_new_protected:Npn \__seq_set_map_x:NNNn #1#2#3#4
8069 {
8070     \__seq_push_item_def:n { \exp_not:N \__seq_item:n {#4} }
8071     #1 #2 { #3 }
8072     \__seq_pop_item_def:
8073 }

```

(End definition for `\seq_set_map_x:NNn`, `\seq_gset_map_x:NNn`, and `__seq_set_map_x:NNNn`. These functions are documented on page 83.)

`\seq_set_map:NNn`
`\seq_gset_map:NNn`
`__seq_set_map:NNNn`

Similar to `\seq_set_map_x:NNn`, but prevents expansion of the `<inline function>`.

```

8074 \cs_new_protected:Npn \seq_set_map:NNn
8075 { \__seq_set_map:NNNn \__kernel_tl_set:Nx }
8076 \cs_new_protected:Npn \seq_gset_map:NNn

```

```

8077 { \__seq_set_map:NNNn \__kernel_tl_gset:Nx }
8078 \cs_new_protected:Npn \__seq_set_map:NNNn #1#2#3#4
8079 {
8080   \__seq_push_item_def:n { \exp_not:n { \__seq_item:n {#4} } }
8081   #1 #2 { #3 }
8082   \__seq_pop_item_def:
8083 }

```

(End definition for `\seq_set_map:NNn`, `\seq_gset_map:NNn`, and `__seq_set_map:NNNn`. These functions are documented on page 82.)

`\seq_count:N` Since counting the items in a sequence is quite common, we optimize it by grabbing 8 items at a time and correspondingly adding 8 to an integer expression. At the end of the loop, #9 is `__seq_count_end:w` instead of being empty. It removes 8+ and instead places the number of `__seq_item:n` that `__seq_count:w` grabbed before reaching the end of the sequence.

```

8084 \cs_new:Npn \seq_count:N #1
8085 {
8086   \int_eval:n
8087   {
8088     \exp_after:wN \use_i:nn
8089     \exp_after:wN \__seq_count:w
8090     #1
8091     \__seq_count_end:w \__seq_item:n 7
8092     \__seq_count_end:w \__seq_item:n 6
8093     \__seq_count_end:w \__seq_item:n 5
8094     \__seq_count_end:w \__seq_item:n 4
8095     \__seq_count_end:w \__seq_item:n 3
8096     \__seq_count_end:w \__seq_item:n 2
8097     \__seq_count_end:w \__seq_item:n 1
8098     \__seq_count_end:w \__seq_item:n 0
8099     \prg_break_point:
8100   }
8101 }
8102 \cs_new:Npn \__seq_count:w
8103   #1 \__seq_item:n #2 \__seq_item:n #3 \__seq_item:n #4 \__seq_item:n
8104   #5 \__seq_item:n #6 \__seq_item:n #7 \__seq_item:n #8 #9 \__seq_item:n
8105   { #9 8 + \__seq_count:w }
8106 \cs_new:Npn \__seq_count_end:w 8 + \__seq_count:w #1#2 \prg_break_point: {#1}
8107 \cs_generate_variant:Nn \seq_count:N { c }

```

(End definition for `\seq_count:N`, `__seq_count:w`, and `__seq_count_end:w`. This function is documented on page 83.)

11.7 Using sequences

`\seq_use:Nnnn` See `\clist_use:Nnnn` for a general explanation. The main difference is that we use `__seq_item:n` as a delimiter rather than commas. We also need to add `__seq_item:n` at various places, and `\s__seq`.

```

\seq_use:cnnn
\__seq_use:NNnNnn
\__seq_use_setup:w
\__seq_use:nwwwnwn
\__seq_use:nwnn
\seq_use:Nn
\seq_use:cn

```

```

8108 \cs_new:Npn \seq_use:Nnnn #1#2#3#4
8109 {
8110   \seq_if_exist:NTF #1
8111   {
8112     \int_case:nnF { \seq_count:N #1 }

```



```

8113     {
8114         { 0 } { }
8115         { 1 } { \exp_after:wN \__seq_use:NnnNnn #1 ? { } { } }
8116         { 2 } { \exp_after:wN \__seq_use:NnnNnn #1 {#2} }
8117     }
8118     {
8119         \exp_after:wN \__seq_use_setup:w #1 \__seq_item:n
8120         \s__seq_mark { \__seq_use:nwwwnwn {#3} }
8121         \s__seq_mark { \__seq_use:nwn {#4} }
8122         \s__seq_stop { }
8123     }
8124 }
8125 {
8126     \__kernel_msg_expandable_error:nnn
8127     { kernel } { bad-variable } {#1}
8128 }
8129 }
8130 \cs_generate_variant:Nn \seq_use:Nnnn { c }
8131 \cs_new:Npn \__seq_use:NnnNnn #1#2#3#4#5#6 { \exp_not:n { #3 #6 #5 } }
8132 \cs_new:Npn \__seq_use_setup:w \s__seq { \__seq_use:nwwwnwn { } }
8133 \cs_new:Npn \__seq_use:nwwwnwn
8134     #1 \__seq_item:n #2 \__seq_item:n #3 \__seq_item:n #4#5
8135     \s__seq_mark #6#7 \s__seq_stop #8
8136     {
8137         #6 \__seq_item:n {#3} \__seq_item:n {#4} #5
8138         \s__seq_mark {#6} #7 \s__seq_stop { #8 #1 #2 }
8139     }
8140 \cs_new:Npn \__seq_use:nwn #1 \__seq_item:n #2 #3 \s__seq_stop #4
8141     { \exp_not:n { #4 #1 #2 } }
8142 \cs_new:Npn \seq_use:Nn #1#2
8143     { \seq_use:Nnnn #1 {#2} {#2} {#2} }
8144 \cs_generate_variant:Nn \seq_use:Nn { c }

```

(End definition for `\seq_use:Nnnn` and others. These functions are documented on page 83.)

11.8 Sequence stacks

The same functions as for sequences, but with the correct naming.

\seq_push:Nn	Pushing to a sequence is the same as adding on the left.
\seq_push:NV	8145 \cs_new_eq:NN \seq_push:Nn \seq_put_left:Nn
\seq_push:Nv	8146 \cs_new_eq:NN \seq_push:Nv \seq_put_left:Nv
\seq_push:No	8147 \cs_new_eq:NN \seq_push:No \seq_put_left:No
\seq_push:Nx	8148 \cs_new_eq:NN \seq_push:Nx \seq_put_left:Nx
\seq_push:cn	8149 \cs_new_eq:NN \seq_push:cn \seq_put_left:cn
\seq_push:cV	8150 \cs_new_eq:NN \seq_push:cV \seq_put_left:cV
\seq_push:cV	8151 \cs_new_eq:NN \seq_push:cV \seq_put_left:cV
\seq_push:co	8152 \cs_new_eq:NN \seq_push:co \seq_put_left:co
\seq_push:cx	8153 \cs_new_eq:NN \seq_push:cx \seq_put_left:cx
\seq_gpush:Nn	8154 \cs_new_eq:NN \seq_gpush:Nn \seq_gput_left:Nn
\seq_gpush:NV	8155 \cs_new_eq:NN \seq_gpush:Nv \seq_gput_left:Nv
\seq_gpush:Nv	8156 \cs_new_eq:NN \seq_gpush:No \seq_gput_left:No
\seq_gpush:No	8157 \cs_new_eq:NN \seq_gpush:Nx \seq_gput_left:Nx
\seq_gpush:Nx	8158 \cs_new_eq:NN \seq_gpush:cn \seq_gput_left:cn
\seq_gpush:cn	
\seq_gpush:cV	
\seq_gpush:cV	
\seq_gpush:co	
\seq_gpush:cx	

```

8159 \cs_new_eq:NN \seq_gpush:Nx \seq_gput_left:Nx
8160 \cs_new_eq:NN \seq_gpush:cn \seq_gput_left:cn
8161 \cs_new_eq:NN \seq_gpush:cV \seq_gput_left:cV
8162 \cs_new_eq:NN \seq_gpush:cv \seq_gput_left:cv
8163 \cs_new_eq:NN \seq_gpush:co \seq_gput_left:co
8164 \cs_new_eq:NN \seq_gpush:cx \seq_gput_left:cx

```

(End definition for `\seq_push:Nn` and `\seq_gpush:Nn`. These functions are documented on page 85.)

`\seq_get:NN` In most cases, getting items from the stack does not need to specify that this is from the left. So alias are provided.

```

\seq_get:cn
\seq_pop:NN
\seq_pop:cn
\seq_gpop:NN
\seq_gpop:cn
8165 \cs_new_eq:NN \seq_get:NN \seq_get_left:NN
8166 \cs_new_eq:NN \seq_get:cn \seq_get_left:cn
8167 \cs_new_eq:NN \seq_pop:NN \seq_pop_left:NN
8168 \cs_new_eq:NN \seq_pop:cn \seq_pop_left:cn
8169 \cs_new_eq:NN \seq_gpop:NN \seq_gpop_left:NN
8170 \cs_new_eq:NN \seq_gpop:cn \seq_gpop_left:cn

```

(End definition for `\seq_get:NN`, `\seq_pop:NN`, and `\seq_gpop:NN`. These functions are documented on page 84.)

`\seq_get:NNTF` More copies.

```

\seq_get:cnTF
\seq_pop:NNTF
\seq_pop:cnTF
\seq_gpop:NNTF
\seq_gpop:cnTF
8171 \prg_new_eq_conditional:NNn \seq_get:NN \seq_get_left:NN { T , F , TF }
8172 \prg_new_eq_conditional:NNn \seq_get:cn \seq_get_left:cn { T , F , TF }
8173 \prg_new_eq_conditional:NNn \seq_pop:NN \seq_pop_left:NN { T , F , TF }
8174 \prg_new_eq_conditional:NNn \seq_pop:cn \seq_pop_left:cn { T , F , TF }
8175 \prg_new_eq_conditional:NNn \seq_gpop:NN \seq_gpop_left:NN { T , F , TF }
8176 \prg_new_eq_conditional:NNn \seq_gpop:cn \seq_gpop_left:cn { T , F , TF }

```

(End definition for `\seq_get:NNTF`, `\seq_pop:NNTF`, and `\seq_gpop:NNTF`. These functions are documented on page 84.)

11.9 Viewing sequences

`\seq_show:N` Apply the general `\msg_show:nnnnnn`.

```

\seq_show:c
\seq_log:N
\seq_log:c
\__seq_show:NN
8177 \cs_new_protected:Npn \seq_show:N { \__seq_show:NN \msg_show:nnxxxx }
8178 \cs_generate_variant:Nn \seq_show:N { c }
8179 \cs_new_protected:Npn \seq_log:N { \__seq_show:NN \msg_log:nnxxxx }
8180 \cs_generate_variant:Nn \seq_log:N { c }
8181 \cs_new_protected:Npn \__seq_show:NN #1#2
8182 {
8183   \__kernel_chk_defined:NT #2
8184   {
8185     #1 { LaTeX/kernel } { show-seq }
8186     { \token_to_str:N #2 }
8187     { \seq_map_function:NN #2 \msg_show_item:n }
8188     { } { }
8189   }
8190 }

```

(End definition for `\seq_show:N`, `\seq_log:N`, and `__seq_show:NN`. These functions are documented on page 87.)

11.10 Scratch sequences

`\l_tmpa_seq` Temporary comma list variables.
`\l_tmpb_seq`
`\g_tmpa_seq`
`\g_tmpb_seq`

```
8191 \seq_new:N \l_tmpa_seq
8192 \seq_new:N \l_tmpb_seq
8193 \seq_new:N \g_tmpa_seq
8194 \seq_new:N \g_tmpb_seq
```

(End definition for `\l_tmpa_seq` and others. These variables are documented on page 87.)

```
8195 \</package>
```

12 l3int implementation

```
8196 \*package>
```

```
8197 \@@=int
```

The following test files are used for this code: `m3int001,m3int002,m3int03`.

`\c_max_register_int` Done in l3basics.

(End definition for `\c_max_register_int`. This variable is documented on page 100.)

`__int_to_roman:w` Done in l3basics.

`\if_int_compare:w` (End definition for `__int_to_roman:w` and `\if_int_compare:w`. This function is documented on page 101.)

`\or:` Done in l3basics.

(End definition for `\or:`. This function is documented on page 101.)

`\int_value:w` Here are the remaining primitives for number comparisons and expressions.

```
\__int_eval:w      8198 \cs_new_eq:NN \int_value:w      \tex_number:D
\__int_eval_end:    8199 \cs_new_eq:NN \__int_eval:w      \tex_numexpr:D
\if_int_odd:w       8200 \cs_new_eq:NN \__int_eval_end:    \tex_relax:D
\if_case:w          8201 \cs_new_eq:NN \if_int_odd:w      \tex_ifodd:D
                    8202 \cs_new_eq:NN \if_case:w        \tex_ifcase:D
```

(End definition for `\int_value:w` and others. These functions are documented on page 101.)

`\s__int_mark` Scan marks used throughout the module.

```
\s__int_stop      8203 \scan_new:N \s__int_mark
                  8204 \scan_new:N \s__int_stop
```

(End definition for `\s__int_mark` and `\s__int_stop`.)

`__int_use_none_delimit_by_s_stop:w` Function to gobble until a scan mark.

```
8205 \cs_new:Npn \__int_use_none_delimit_by_s_stop:w #1 \s__int_stop { }
```

(End definition for `__int_use_none_delimit_by_s_stop:w`.)

`\q__int_recursion_tail` Quarks for recursion.

```
\q__int_recursion_stop 8206 \quark_new:N \q__int_recursion_tail
                        8207 \quark_new:N \q__int_recursion_stop
```

(End definition for `\q__int_recursion_tail` and `\q__int_recursion_stop`.)

_int_if_recursion_tail_stop_do:Nn
_int_if_recursion_tail_stop:N

Functions to query quarks.

```
8208 \_kernel_quark_new_test:N \_int_if_recursion_tail_stop_do:Nn
8209 \_kernel_quark_new_test:N \_int_if_recursion_tail_stop:N
```

(End definition for _int_if_recursion_tail_stop_do:Nn and _int_if_recursion_tail_stop:N.)

12.1 Integer expressions

\int_eval:n Wrapper for _int_eval:w: can be used in an integer expression or directly in the input stream. When debugging, use parentheses to catch early termination.

\int_eval:w

```
8210 \cs_new:Npn \int_eval:n #1
8211 { \int_value:w \_int_eval:w #1 \_int_eval_end: }
8212 \cs_new:Npn \int_eval:w { \int_value:w \_int_eval:w }
```

(End definition for \int_eval:n and \int_eval:w. These functions are documented on page 89.)

\int_sign:n
_int_sign:Nw

See \int_abs:n. Evaluate the expression once (and when debugging is enabled, check that the expression is well-formed), then test the first character to determine the sign. This is wrapped in \int_value:w ... \exp_stop_f: to ensure a fixed number of expansions and to avoid dealing with closing the conditionals.

```
8213 \cs_new:Npn \int_sign:n #1
8214 {
8215   \int_value:w \exp_after:wN \_int_sign:Nw
8216   \int_value:w \_int_eval:w #1 \_int_eval_end: ;
8217   \exp_stop_f:
8218 }
8219 \cs_new:Npn \_int_sign:Nw #1#2 ;
8220 {
8221   \if_meaning:w 0 #1
8222   0
8223   \else:
8224     \if_meaning:w - #1 - \fi: 1
8225   \fi:
8226 }
```

(End definition for \int_sign:n and _int_sign:Nw. This function is documented on page 90.)

\int_abs:n
_int_abs:N

Functions for min, max, and absolute value with only one evaluation. The absolute value is obtained by removing a leading sign if any. All three functions expand in two steps.

\int_max:nn

\int_min:nn

_int_maxmin:wwN

```
8227 \cs_new:Npn \int_abs:n #1
8228 {
8229   \int_value:w \exp_after:wN \_int_abs:N
8230   \int_value:w \_int_eval:w #1 \_int_eval_end:
8231   \exp_stop_f:
8232 }
8233 \cs_new:Npn \_int_abs:N #1
8234 { \if_meaning:w - #1 \else: \exp_after:wN #1 \fi: }
8235 \cs_set:Npn \int_max:nn #1#2
8236 {
8237   \int_value:w \exp_after:wN \_int_maxmin:wwN
8238   \int_value:w \_int_eval:w #1 \exp_after:wN ;
8239   \int_value:w \_int_eval:w #2 ;
8240   >
8241   \exp_stop_f:
```

```

8242 }
8243 \cs_set:Npn \int_min:nn #1#2
8244 {
8245   \int_value:w \exp_after:wN \__int_maxmin:wwN
8246   \int_value:w \__int_eval:w #1 \exp_after:wN ;
8247   \int_value:w \__int_eval:w #2 ;
8248   <
8249   \exp_stop_f:
8250 }
8251 \cs_new:Npn \__int_maxmin:wwN #1 ; #2 ; #3
8252 {
8253   \if_int_compare:w #1 #3 #2 ~
8254     #1
8255   \else:
8256     #2
8257   \fi:
8258 }

```

(End definition for `\int_abs:n` and others. These functions are documented on page 90.)

`\int_div_truncate:nn` As `__int_eval:w` rounds the result of a division we also provide a version that truncates the result. We use an auxiliary to make sure numerator and denominator are only evaluated once: this comes in handy when those are more expressions are expensive to evaluate (e.g., `\tl_count:n`). If the numerator `#1#2` is 0, then we divide 0 by the denominator (this ensures that 0/0 is correctly reported as an error). Otherwise, shift the numerator `#1#2` towards 0 by $(| \#3\#4 | - 1)/2$, which we round away from zero. It turns out that this quantity exactly compensates the difference between ε -TeX's rounding and the truncating behaviour that we want. The details are thanks to Heiko Oberdiek: getting things right in all cases is not so easy.

```

8259 \cs_new:Npn \int_div_truncate:nn #1#2
8260 {
8261   \int_value:w \__int_eval:w
8262   \exp_after:wN \__int_div_truncate:NwNw
8263   \int_value:w \__int_eval:w #1 \exp_after:wN ;
8264   \int_value:w \__int_eval:w #2 ;
8265   \__int_eval_end:
8266 }
8267 \cs_new:Npn \__int_div_truncate:NwNw #1#2; #3#4;
8268 {
8269   \if_meaning:w 0 #1
8270     0
8271   \else:
8272     (
8273       #1#2
8274       \if_meaning:w - #1 + \else: - \fi:
8275       ( \if_meaning:w - #3 - \fi: #3#4 - 1 ) / 2
8276     )
8277   \fi:
8278   / #3#4
8279 }

```

For the sake of completeness:

```

8280 \cs_new:Npn \int_div_round:nn #1#2
8281 { \int_value:w \__int_eval:w ( #1 ) / ( #2 ) \__int_eval_end: }

```

Finally there's the modulus operation.

```

8282 \cs_new:Npn \int_mod:nn #1#2
8283 {
8284   \int_value:w \__int_eval:w \exp_after:wN \__int_mod:ww
8285   \int_value:w \__int_eval:w #1 \exp_after:wN ;
8286   \int_value:w \__int_eval:w #2 ;
8287   \__int_eval_end:
8288 }
8289 \cs_new:Npn \__int_mod:ww #1; #2;
8290 { #1 - ( \__int_div_truncate:NwNw #1 ; #2 ; ) * #2 }

```

(End definition for `\int_div_truncate:nn` and others. These functions are documented on page 90.)

`__kernel_int_add:nnn` Equivalent to `\int_eval:n {#1+#2+#3}` except that overflow only occurs if the final result overflows $[-2^{31} + 1, 2^{31} - 1]$. The idea is to choose the order in which the three numbers are added together. If #1 and #2 have opposite signs (one is in $[-2^{31} + 1, -1]$ and the other in $[0, 2^{31} - 1]$) then #1+#2 cannot overflow so we compute the result as #1+#2+#3. If they have the same sign, then either #3 has the same sign and the order does not matter, or #3 has the opposite sign and any order in which #3 is not last will work. We use #1+#3+#2.

```

8291 \cs_new:Npn \__kernel_int_add:nnn #1#2#3
8292 {
8293   \int_value:w \__int_eval:w #1
8294   \if_int_compare:w #2 < \c_zero_int \exp_after:wN \reverse_if:N \fi:
8295   \if_int_compare:w #1 < \c_zero_int + #2 + #3 \else: + #3 + #2 \fi:
8296   \__int_eval_end:
8297 }

```

(End definition for `__kernel_int_add:nnn`.)

12.2 Creating and initialising integers

`\int_new:N` Two ways to do this: one for the format and one for the L^AT_EX 2_ε package. In plain T_EX, `\int_new:c` `\newcount` (and other allocators) are `\outer:` to allow the code here to work in “generic” mode this is therefore accessed by name. (The same applies to `\newbox`, `\newdimen` and so on.)

```

8298 \cs_new_protected:Npn \int_new:N #1
8299 {
8300   \__kernel_chk_if_free_cs:N #1
8301   \cs:w newcount \cs_end: #1
8302 }
8303 \cs_generate_variant:Nn \int_new:N { c }

```

(End definition for `\int_new:N`. This function is documented on page 90.)

`\int_const:Nn` As stated, most constants can be defined as `\chardef` or `\mathchardef` but that's engine dependent. As a result, there is some set up code to determine what can be done. No full engine testing just yet so everything is a little awkward. We cannot use `\int_gset:Nn` because (when `check-declarations` is enabled) this runs some checks that constants would fail.

```

\int_const:cn
\__int_constdef:Nw
\c__int_max_constdef_int
8304 \cs_new_protected:Npn \int_const:Nn #1#2
8305 {
8306   \int_compare:nNnTF {#2} < \c_zero_int

```

```

8307     {
8308         \int_new:N #1
8309         \tex_global:D
8310     }
8311     {
8312         \int_compare:nNnTF {#2} > \c__int_max_constdef_int
8313         {
8314             \int_new:N #1
8315             \tex_global:D
8316         }
8317         {
8318             \__kernel_chk_if_free_cs:N #1
8319             \tex_global:D \__int_constdef:Nw
8320         }
8321     }
8322     #1 = \__int_eval:w #2 \__int_eval_end:
8323 }
8324 \cs_generate_variant:Nn \int_const:Nn { c }
8325 \if_int_odd:w 0
8326     \cs_if_exist:NT \tex_luatexversion:D { 1 }
8327     \cs_if_exist:NT \tex_omathchardef:D { 1 }
8328     \cs_if_exist:NT \tex_XeTeXversion:D { 1 } ~
8329     \cs_if_exist:NTF \tex_omathchardef:D
8330     { \cs_new_eq:NN \__int_constdef:Nw \tex_omathchardef:D }
8331     { \cs_new_eq:NN \__int_constdef:Nw \tex_chardef:D }
8332     \__int_constdef:Nw \c__int_max_constdef_int 1114111 ~
8333 \else:
8334     \cs_new_eq:NN \__int_constdef:Nw \tex_mathchardef:D
8335     \tex_mathchardef:D \c__int_max_constdef_int 32767 ~
8336 \fi:

```

(End definition for `\int_const:Nn`, `__int_constdef:Nw`, and `\c__int_max_constdef_int`. This function is documented on page 91.)

`\int_zero:N` Functions that reset an *integer* register to zero.

```

\int_zero:c 8337 \cs_new_protected:Npn \int_zero:N #1 { #1 = \c_zero_int }
\int_gzero:N 8338 \cs_new_protected:Npn \int_gzero:N #1 { \tex_global:D #1 = \c_zero_int }
\int_gzero:c 8339 \cs_generate_variant:Nn \int_zero:N { c }
8340 \cs_generate_variant:Nn \int_gzero:N { c }

```

(End definition for `\int_zero:N` and `\int_gzero:N`. These functions are documented on page 91.)

`\int_zero_new:N` Create a register if needed, otherwise clear it.

```

\int_zero_new:c 8341 \cs_new_protected:Npn \int_zero_new:N #1
\int_gzero_new:N 8342 { \int_if_exist:NTF #1 { \int_zero:N #1 } { \int_new:N #1 } }
\int_gzero_new:c 8343 \cs_new_protected:Npn \int_gzero_new:N #1
8344 { \int_if_exist:NTF #1 { \int_gzero:N #1 } { \int_new:N #1 } }
8345 \cs_generate_variant:Nn \int_zero_new:N { c }
8346 \cs_generate_variant:Nn \int_gzero_new:N { c }

```

(End definition for `\int_zero_new:N` and `\int_gzero_new:N`. These functions are documented on page 91.)

`\int_set_eq:NN` Setting equal means using one integer inside the set function of another. Check that assigned integer is local/global. No need to check that the other one is defined as `\TeX` does it for us.

`\int_set_eq:cN`

`\int_set_eq:Nc`

`\int_set_eq:cc` 8347 `\cs_new_protected:Npn \int_set_eq:NN #1#2 { #1 = #2 }`

`\int_gset_eq:NN` 8348 `\cs_generate_variant:Nn \int_set_eq:NN { c , Nc , cc }`

`\int_gset_eq:cN` 8349 `\cs_new_protected:Npn \int_gset_eq:NN #1#2 { \tex_global:D #1 = #2 }`

`\int_gset_eq:Nc` 8350 `\cs_generate_variant:Nn \int_gset_eq:NN { c , Nc , cc }`

`\int_gset_eq:cc`

(End definition for `\int_set_eq:NN` and `\int_gset_eq:NN`. These functions are documented on page 91.)

`\int_if_exist:p:N` Copies of the `cs` functions defined in `l3basics`.

`\int_if_exist:p:c` 8351 `\prg_new_eq_conditional:NNn \int_if_exist:N \cs_if_exist:N`

`\int_if_exist:NTF` 8352 `{ TF , T , F , p }`

`\int_if_exist:cTF` 8353 `\prg_new_eq_conditional:NNn \int_if_exist:c \cs_if_exist:c`

8354 `{ TF , T , F , p }`

(End definition for `\int_if_exist:N`. This function is documented on page 91.)

12.3 Setting and incrementing integers

`\int_add:Nn` Adding and subtracting to and from a counter.

`\int_add:cn` 8355 `\cs_new_protected:Npn \int_add:Nn #1#2`

`\int_gadd:Nn` 8356 `{ \tex_advance:D #1 by __int_eval:w #2 __int_eval_end: }`

`\int_gadd:cn` 8357 `\cs_new_protected:Npn \int_sub:Nn #1#2`

`\int_sub:Nn` 8358 `{ \tex_advance:D #1 by - __int_eval:w #2 __int_eval_end: }`

`\int_sub:cn` 8359 `\cs_new_protected:Npn \int_gadd:Nn #1#2`

`\int_gsub:Nn` 8360 `{ \tex_global:D \tex_advance:D #1 by __int_eval:w #2 __int_eval_end: }`

`\int_gsub:cn` 8361 `\cs_new_protected:Npn \int_gsub:Nn #1#2`

8362 `{ \tex_global:D \tex_advance:D #1 by - __int_eval:w #2 __int_eval_end: }`

8363 `\cs_generate_variant:Nn \int_add:Nn { c }`

8364 `\cs_generate_variant:Nn \int_gadd:Nn { c }`

8365 `\cs_generate_variant:Nn \int_sub:Nn { c }`

8366 `\cs_generate_variant:Nn \int_gsub:Nn { c }`

(End definition for `\int_add:Nn` and others. These functions are documented on page 91.)

`\int_incr:N` Incrementing and decrementing of integer registers is done with the following functions.

`\int_incr:c` 8367 `\cs_new_protected:Npn \int_incr:N #1`

`\int_gincr:N` 8368 `{ \tex_advance:D #1 \c_one_int }`

`\int_gincr:c` 8369 `\cs_new_protected:Npn \int_decr:N #1`

`\int_decr:N` 8370 `{ \tex_advance:D #1 - \c_one_int }`

`\int_decr:c` 8371 `\cs_new_protected:Npn \int_gincr:N #1`

`\int_gdecr:N` 8372 `{ \tex_global:D \tex_advance:D #1 \c_one_int }`

`\int_gdecr:c` 8373 `\cs_new_protected:Npn \int_gdecr:N #1`

8374 `{ \tex_global:D \tex_advance:D #1 - \c_one_int }`

8375 `\cs_generate_variant:Nn \int_incr:N { c }`

8376 `\cs_generate_variant:Nn \int_decr:N { c }`

8377 `\cs_generate_variant:Nn \int_gincr:N { c }`

8378 `\cs_generate_variant:Nn \int_gdecr:N { c }`

(End definition for `\int_incr:N` and others. These functions are documented on page 91.)

`\int_set:Nn` As integers are register-based T_EX issues an error if they are not defined.

```

\int_set:cn      8379 \cs_new_protected:Npn \int_set:Nn #1#2
\int_gset:Nn      8380 { #1 ~ \__int_eval:w #2 \__int_eval_end: }
\int_gset:cn      8381 \cs_new_protected:Npn \int_gset:Nn #1#2
                  8382 { \tex_global:D #1 ~ \__int_eval:w #2 \__int_eval_end: }
                  8383 \cs_generate_variant:Nn \int_set:Nn { c }
                  8384 \cs_generate_variant:Nn \int_gset:Nn { c }

```

(End definition for `\int_set:Nn` and `\int_gset:Nn`. These functions are documented on page 92.)

12.4 Using integers

`\int_use:N` Here is how counters are accessed:

```

\int_use:c      8385 \cs_new_eq:NN \int_use:N \tex_the:D
We hand-code this for some speed gain:
8386 %\cs_generate_variant:Nn \int_use:N { c }
8387 \cs_new:Npn \int_use:c #1 { \tex_the:D \cs:w #1 \cs_end: }

```

(End definition for `\int_use:N`. This function is documented on page 92.)

12.5 Integer expression conditionals

`__int_compare_error:` Those functions are used for comparison tests which use a simple syntax where only one set of braces is required and additional operators such as `!=` and `>=` are supported. `__int_compare_error:Nw` The tests first evaluate their left-hand side, with a trailing `__int_compare_error:.` This marker is normally not expanded, but if the relation symbol is missing from the test's argument, then the marker inserts `=` (and itself) after triggering the relevant T_EX error. If the first token which appears after evaluating and removing the left-hand side is not a known relation symbol, then a judiciously placed `__int_compare_error:Nw` gets expanded, cleaning up the end of the test and telling the user what the problem was.

```

8388 \cs_new_protected:Npn \__int_compare_error:
8389 {
8390   \if_int_compare:w \c_zero_int \c_zero_int \fi:
8391   =
8392   \__int_compare_error:
8393 }
8394 \cs_new:Npn \__int_compare_error:Nw
8395 #1#2 \s__int_stop
8396 {
8397   { }
8398   \c_zero_int \fi:
8399   \__kernel_msg_expandable_error:nnn
8400   { kernel } { unknown-comparison } {#1}
8401   \prg_return_false:
8402 }

```

(End definition for `__int_compare_error:` and `__int_compare_error:Nw`.)

`\int_compare_p:n` Comparison tests using a simple syntax where only one set of braces is required, additional operators such as `!=` and `>=` are supported, and multiple comparisons can be performed at once, for instance `0 < 5 <= 1`. The idea is to loop through the argument, finding one operand at a time, and comparing it to the previous one. The looping auxiliary `__int_compare:Nw` reads one *operand* and one *comparison* symbol, and leaves roughly

```

\__int_compare_end=:NNw
\__int_compare=:NNw
\__int_compare<:NNw
\__int_compare>:NNw
\__int_compare==:NNw
\__int_compare!=:NNw
\__int_compare<=:NNw
\__int_compare>=:NNw

```

```

\operand\prg_return_false:\fi:
\reverse_if:N\if_int_compare:w\operand\comparison
\__int_compare:Nw

```

in the input stream. Each call to this auxiliary provides the second operand of the last call's `\if_int_compare:w`. If one of the *comparisons* is `false`, the `true` branch of the \TeX conditional is taken (because of `\reverse_if:N`), immediately returning `false` as the result of the test. There is no \TeX conditional waiting the first operand, so we add an `\if_false:` and expand by hand with `\int_value:w`, thus skipping `\prg_return_false:` on the first iteration.

Before starting the loop, the first step is to make sure that there is at least one relation symbol. We first let \TeX evaluate this left hand side of the (in)equality using `__int_eval:w`. Since the relation symbols `<`, `>`, `=` and `!` are not allowed in integer expressions, they would terminate the expression. If the argument contains no relation symbol, `__int_compare_error:` is expanded, inserting `=` and itself after an error. In all cases, `__int_compare:w` receives as its argument an integer, a relation symbol, and some more tokens. We then setup the loop, which is ended by the two odd-looking items `e` and `{=nd_}`, with a trailing `\s__int_stop` used to grab the entire argument when necessary.

```

8403 \prg_new_conditional:Npnn \int_compare:n #1 { p , T , F , TF }
8404 {
8405   \exp_after:wN \__int_compare:w
8406   \int_value:w \__int_eval:w #1 \__int_compare_error:
8407 }
8408 \cs_new:Npn \__int_compare:w #1 \__int_compare_error:
8409 {
8410   \exp_after:wN \if_false: \int_value:w
8411   \__int_compare:Nw #1 e { = nd_ } \s__int_stop
8412 }

```

The goal here is to find an *operand* and a *comparison*. The *operand* is already evaluated, but we cannot yet grab it as an argument. To access the following relation symbol, we remove the number by applying `__int_to_roman:w`, after making sure that the argument becomes non-positive: its roman numeral representation is then empty. Then probe the first two tokens with `__int_compare:NNw` to determine the relation symbol, building a control sequence from it (`\token_to_str:N` gives better errors if `#1` is not a character). All the extended forms have an extra `=` hence the test for that as a second token. If the relation symbol is unknown, then the control sequence is turned by \TeX into `\scan_stop:`, ignored thanks to `\unexpanded`, and `__int_compare_error:Nw` raises an error.

```

8413 \cs_new:Npn \__int_compare:Nw #1#2 \s__int_stop
8414 {
8415   \exp_after:wN \__int_compare:NNw
8416   \__int_to_roman:w - 0 #2 \s__int_mark
8417   #1#2 \s__int_stop
8418 }
8419 \cs_new:Npn \__int_compare:NNw #1#2#3 \s__int_mark
8420 {
8421   \__kernel_exp_not:w
8422   \use:c
8423   {
8424     __int_compare_ \token_to_str:N #1

```

```

8425         \if_meaning:w = #2 = \fi:
8426         :NNw
8427     }
8428     \__int_compare_error:Nw #1
8429 }

```

When the last *operand* is seen, `__int_compare:NNw` receives `e` and `=nd_` as arguments, hence calling `__int_compare_end=:NNw` to end the loop: return the result of the last comparison (involving the operand that we just found). When a normal relation is found, the appropriate auxiliary calls `__int_compare:nnN` where `#1` is `\if_int_compare:w` or `\reverse_if:N \if_int_compare:w`, `#2` is the *operand*, and `#3` is one of `<`, `=`, or `>`. As announced earlier, we leave the *operand* for the previous conditional. If this conditional is true the result of the test is known, so we remove all tokens and return `false`. Otherwise, we apply the conditional `#1` to the *operand* `#2` and the comparison `#3`, and call `__int_compare:Nw` to look for additional operands, after evaluating the following expression.

```

8430 \cs_new:cpn { __int_compare_end=:NNw } #1#2#3 e #4 \s__int_stop
8431 {
8432     {#3} \exp_stop_f:
8433     \prg_return_false: \else: \prg_return_true: \fi:
8434 }
8435 \cs_new:Npn \__int_compare:nnN #1#2#3
8436 {
8437     {#2} \exp_stop_f:
8438     \prg_return_false: \exp_after:wN \__int_use_none_delimit_by_s_stop:w
8439     \fi:
8440     #1 #2 #3 \exp_after:wN \__int_compare:Nw \int_value:w \__int_eval:w
8441 }

```

The actual comparisons are then simple function calls, using the relation as delimiter for a delimited argument and discarding `__int_compare_error:Nw` *token* responsible for error detection.

```

8442 \cs_new:cpn { __int_compare=:NNw } #1#2#3 =
8443 { \__int_compare:nnN { \reverse_if:N \if_int_compare:w } {#3} = }
8444 \cs_new:cpn { __int_compare:<:NNw } #1#2#3 <
8445 { \__int_compare:nnN { \reverse_if:N \if_int_compare:w } {#3} < }
8446 \cs_new:cpn { __int_compare:>:NNw } #1#2#3 >
8447 { \__int_compare:nnN { \reverse_if:N \if_int_compare:w } {#3} > }
8448 \cs_new:cpn { __int_compare=:NNw } #1#2#3 ==
8449 { \__int_compare:nnN { \reverse_if:N \if_int_compare:w } {#3} = }
8450 \cs_new:cpn { __int_compare!=:NNw } #1#2#3 !=
8451 { \__int_compare:nnN { \if_int_compare:w } {#3} = }
8452 \cs_new:cpn { __int_compare<=:NNw } #1#2#3 <=
8453 { \__int_compare:nnN { \if_int_compare:w } {#3} > }
8454 \cs_new:cpn { __int_compare_>=:NNw } #1#2#3 >=
8455 { \__int_compare:nnN { \if_int_compare:w } {#3} < }

```

(End definition for `\int_compare:nTF` and others. This function is documented on page 93.)

`\int_compare_p:nNn`
`\int_compare:nNnTF`

More efficient but less natural in typing.

```

8456 \prg_new_conditional:Npnn \int_compare:nNn #1#2#3 { p , T , F , TF }
8457 {
8458     \if_int_compare:w \__int_eval:w #1 #2 \__int_eval:w #3 \__int_eval_end:
8459     \prg_return_true:

```

```

8460     \else:
8461         \prg_return_false:
8462     \fi:
8463 }

```

(End definition for \int_compare:nNnTF. This function is documented on page 92.)

```

\int_case:nn For integer cases, the first task to fully expand the check condition. The over all idea is
\int_case:nnTF then much the same as for \tl_case:nn(TF) as described in l3tl.
\__int_case:nnTF
\__int_case:nw
\__int_case_end:nw
8464 \cs_new:Npn \int_case:nnTF #1
8465 {
8466     \exp:w
8467     \exp_args:Nf \__int_case:nnTF { \int_eval:n {#1} }
8468 }
8469 \cs_new:Npn \int_case:nnT #1#2#3
8470 {
8471     \exp:w
8472     \exp_args:Nf \__int_case:nnTF { \int_eval:n {#1} } {#2} {#3} { }
8473 }
8474 \cs_new:Npn \int_case:nnF #1#2
8475 {
8476     \exp:w
8477     \exp_args:Nf \__int_case:nnTF { \int_eval:n {#1} } {#2} { }
8478 }
8479 \cs_new:Npn \int_case:nn #1#2
8480 {
8481     \exp:w
8482     \exp_args:Nf \__int_case:nnTF { \int_eval:n {#1} } {#2} { } { }
8483 }
8484 \cs_new:Npn \__int_case:nnTF #1#2#3#4
8485 { \__int_case:nw {#1} #2 {#1} { } \s__int_mark {#3} \s__int_mark {#4} \s__int_stop }
8486 \cs_new:Npn \__int_case:nw #1#2#3
8487 {
8488     \int_compare:nNnTF {#1} = {#2}
8489     { \__int_case_end:nw {#3} }
8490     { \__int_case:nw {#1} }
8491 }
8492 \cs_new:Npn \__int_case_end:nw #1#2#3 \s__int_mark #4#5 \s__int_stop
8493 { \exp_end: #1 #4 }

```

(End definition for \int_case:nnTF and others. This function is documented on page 94.)

```

\int_if_odd_p:n A predicate function.
\int_if_odd:nTF
\int_if_even_p:n
\int_if_even:nTF
8494 \prg_new_conditional:Npnn \int_if_odd:n #1 { p , T , F , TF}
8495 {
8496     \if_int_odd:w \__int_eval:w #1 \__int_eval_end:
8497     \prg_return_true:
8498     \else:
8499     \prg_return_false:
8500     \fi:
8501 }
8502 \prg_new_conditional:Npnn \int_if_even:n #1 { p , T , F , TF}
8503 {
8504     \reverse_if:N \if_int_odd:w \__int_eval:w #1 \__int_eval_end:

```

```

8505     \prg_return_true:
8506 \else:
8507     \prg_return_false:
8508 \fi:
8509 }

```

(End definition for `\int_if_odd:nTF` and `\int_if_even:nTF`. These functions are documented on page 94.)

12.6 Integer expression loops

`\int_while_do:nn` These are quite easy given the above functions. The `while` versions test first and then execute the body. The `do_while` does it the other way round.

```

\int_until_do:nn
\int_do_while:nn
\int_do_until:nn
8510 \cs_new:Npn \int_while_do:nn #1#2
8511 {
8512     \int_compare:nT {#1}
8513     {
8514         #2
8515         \int_while_do:nn {#1} {#2}
8516     }
8517 }
8518 \cs_new:Npn \int_until_do:nn #1#2
8519 {
8520     \int_compare:nF {#1}
8521     {
8522         #2
8523         \int_until_do:nn {#1} {#2}
8524     }
8525 }
8526 \cs_new:Npn \int_do_while:nn #1#2
8527 {
8528     #2
8529     \int_compare:nT {#1}
8530     { \int_do_while:nn {#1} {#2} }
8531 }
8532 \cs_new:Npn \int_do_until:nn #1#2
8533 {
8534     #2
8535     \int_compare:nF {#1}
8536     { \int_do_until:nn {#1} {#2} }
8537 }

```

(End definition for `\int_while_do:nn` and others. These functions are documented on page 95.)

`\int_while_do:nNnn` As above but not using the more natural syntax.

```

\int_until_do:nNnn
\int_do_while:nNnn
\int_do_until:nNnn
8538 \cs_new:Npn \int_while_do:nNnn #1#2#3#4
8539 {
8540     \int_compare:nNnT {#1} #2 {#3}
8541     {
8542         #4
8543         \int_while_do:nNnn {#1} #2 {#3} {#4}
8544     }
8545 }
8546 \cs_new:Npn \int_until_do:nNnn #1#2#3#4

```

```

8547 {
8548   \int_compare:nNnF {#1} #2 {#3}
8549   {
8550     #4
8551     \int_until_do:nNnn {#1} #2 {#3} {#4}
8552   }
8553 }
8554 \cs_new:Npn \int_do_while:nNnn #1#2#3#4
8555 {
8556   #4
8557   \int_compare:nNnT {#1} #2 {#3}
8558   { \int_do_while:nNnn {#1} #2 {#3} {#4} }
8559 }
8560 \cs_new:Npn \int_do_until:nNnn #1#2#3#4
8561 {
8562   #4
8563   \int_compare:nNnF {#1} #2 {#3}
8564   { \int_do_until:nNnn {#1} #2 {#3} {#4} }
8565 }

```

(End definition for `\int_while_do:nNnn` and others. These functions are documented on page 95.)

12.7 Integer step functions

`\int_step_function:nnnN` Before all else, evaluate the initial value, step, and final value. Repeating a function by steps first needs a check on the direction of the steps. After that, do the function for the start value then step and loop around. It would be more symmetrical to test for a step size of zero before checking the sign, but we optimize for the most frequent case (positive step).

```

8566 \cs_new:Npn \int_step_function:nnnN #1#2#3
8567 {
8568   \exp_after:wN \__int_step:wwwN
8569   \int_value:w \__int_eval:w #1 \exp_after:wN ;
8570   \int_value:w \__int_eval:w #2 \exp_after:wN ;
8571   \int_value:w \__int_eval:w #3 ;
8572 }
8573 \cs_new:Npn \__int_step:wwwN #1; #2; #3; #4
8574 {
8575   \int_compare:nNnTF {#2} > \c_zero_int
8576   { \__int_step:NwnnnN > }
8577   {
8578     \int_compare:nNnTF {#2} = \c_zero_int
8579     {
8580       \__kernel_msg_expandable_error:nnn
8581       { kernel } { zero-step } {#4}
8582       \prg_break:
8583     }
8584     { \__int_step:NwnnnN < }
8585   }
8586   #1 ; {#2} {#3} #4
8587   \prg_break_point:
8588 }
8589 \cs_new:Npn \__int_step:NwnnnN #1#2 ; #3#4#5
8590 {

```

```

8591 \if_int_compare:w #2 #1 #4 \exp_stop_f:
8592 \prg_break:n
8593 \fi:
8594 #5 {#2}
8595 \exp_after:wN \__int_step:NwnnN
8596 \exp_after:wN #1
8597 \int_value:w \__int_eval:w #2 + #3 ; {#3} {#4} #5
8598 }
8599 \cs_new:Npn \int_step_function:nN
8600 { \int_step_function:nnnN { 1 } { 1 } }
8601 \cs_new:Npn \int_step_function:nnN #1
8602 { \int_step_function:nnnN {#1} { 1 } }

```

(End definition for `\int_step_function:nnnN` and others. These functions are documented on page 96.)

`\int_step_inline:nn` The approach here is to build a function, with a global integer required to make the nesting safe (as seen in other in line functions), and map that function using `\int_step_function:nnnN`. We put a `\prg_break_point:Nn` so that `map_break` functions from other modules correctly decrement `\g__kernel_prg_map_int` before looking for their own break point. The first argument is `\scan_stop:`, so that no breaking function recognizes this break point as its own.

```

\int_step_inline:nnnN
\int_step_inline:nnnnN
\int_step_variable:nNn
\int_step_variable:nnNn
\int_step_variable:nnnNn
\__int_step:NNnnnn
8603 \cs_new_protected:Npn \int_step_inline:nn
8604 { \int_step_inline:nnnn { 1 } { 1 } }
8605 \cs_new_protected:Npn \int_step_inline:nnn #1
8606 { \int_step_inline:nnnn {#1} { 1 } }
8607 \cs_new_protected:Npn \int_step_inline:nnnn
8608 {
8609 \int_gincr:N \g__kernel_prg_map_int
8610 \exp_args:NNc \__int_step:NNnnnn
8611 \cs_gset_protected:Npn
8612 { \__int_map_ \int_use:N \g__kernel_prg_map_int :w }
8613 }
8614 \cs_new_protected:Npn \int_step_variable:nNn
8615 { \int_step_variable:nnnNn { 1 } { 1 } }
8616 \cs_new_protected:Npn \int_step_variable:nnNn #1
8617 { \int_step_variable:nnnNn {#1} { 1 } }
8618 \cs_new_protected:Npn \int_step_variable:nnnNn #1#2#3#4#5
8619 {
8620 \int_gincr:N \g__kernel_prg_map_int
8621 \exp_args:NNc \__int_step:NNnnnn
8622 \cs_gset_protected:Npx
8623 { \__int_map_ \int_use:N \g__kernel_prg_map_int :w }
8624 {#1}{#2}{#3}
8625 {
8626 \tl_set:Nn \exp_not:N #4 {##1}
8627 \exp_not:n {#5}
8628 }
8629 }
8630 \cs_new_protected:Npn \__int_step:NNnnnn #1#2#3#4#5#6
8631 {
8632 #1 #2 ##1 {#6}
8633 \int_step_function:nnnN {#3} {#4} {#5} #2
8634 \prg_break_point:Nn \scan_stop: { \int_gdecr:N \g__kernel_prg_map_int }
8635 }

```

(End definition for `\int_step_inline:nn` and others. These functions are documented on page 96.)

12.8 Formatting integers

`\int_to_arabic:n` Nothing exciting here.

```
8636 \cs_new_eq:NN \int_to_arabic:n \int_eval:n
```

(End definition for `\int_to_arabic:n`. This function is documented on page 97.)

`\int_to_symbols:nnn` For conversion of integers to arbitrary symbols the method is in general as follows. The input number (#1) is compared to the total number of symbols available at each place (#2). If the input is larger than the total number of symbols available then the modulus is needed, with one added so that the positions don't have to number from zero. Using an `f`-type expansion, this is done so that the system is recursive. The actual conversion function therefore gets a 'nice' number at each stage. Of course, if the initial input was small enough then there is no problem and everything is easy.

```
8637 \cs_new:Npn \int_to_symbols:nnn #1#2#3
8638 {
8639   \int_compare:nNnTF {#1} > {#2}
8640   {
8641     \exp_args:NNo \exp_args:No \__int_to_symbols:nnnn
8642     {
8643       \int_case:nn
8644         { 1 + \int_mod:nn { #1 - 1 } {#2} }
8645         {#3}
8646     }
8647     {#1} {#2} {#3}
8648   }
8649   { \int_case:nn {#1} {#3} }
8650 }
8651 \cs_new:Npn \__int_to_symbols:nnnn #1#2#3#4
8652 {
8653   \exp_args:Nf \int_to_symbols:nnn
8654   { \int_div_truncate:nn { #2 - 1 } {#3} } {#3} {#4}
8655   #1
8656 }
```

(End definition for `\int_to_symbols:nnn` and `__int_to_symbols:nnnn`. This function is documented on page 97.)

`\int_to_alph:n` These both use the above function with input functions that make sense for the alphabet
`\int_to_Alph:n` in English.

```
8657 \cs_new:Npn \int_to_alph:n #1
8658 {
8659   \int_to_symbols:nnn {#1} { 26 }
8660   {
8661     { 1 } { a }
8662     { 2 } { b }
8663     { 3 } { c }
8664     { 4 } { d }
8665     { 5 } { e }
8666     { 6 } { f }
8667     { 7 } { g }
```



```

8668         { 8 } { h }
8669         { 9 } { i }
8670         { 10 } { j }
8671         { 11 } { k }
8672         { 12 } { l }
8673         { 13 } { m }
8674         { 14 } { n }
8675         { 15 } { o }
8676         { 16 } { p }
8677         { 17 } { q }
8678         { 18 } { r }
8679         { 19 } { s }
8680         { 20 } { t }
8681         { 21 } { u }
8682         { 22 } { v }
8683         { 23 } { w }
8684         { 24 } { x }
8685         { 25 } { y }
8686         { 26 } { z }
8687     }
8688 }
8689 \cs_new:Npn \int_to_Alph:n #1
8690 {
8691     \int_to_symbols:nnn {#1} { 26 }
8692     {
8693         { 1 } { A }
8694         { 2 } { B }
8695         { 3 } { C }
8696         { 4 } { D }
8697         { 5 } { E }
8698         { 6 } { F }
8699         { 7 } { G }
8700         { 8 } { H }
8701         { 9 } { I }
8702         { 10 } { J }
8703         { 11 } { K }
8704         { 12 } { L }
8705         { 13 } { M }
8706         { 14 } { N }
8707         { 15 } { O }
8708         { 16 } { P }
8709         { 17 } { Q }
8710         { 18 } { R }
8711         { 19 } { S }
8712         { 20 } { T }
8713         { 21 } { U }
8714         { 22 } { V }
8715         { 23 } { W }
8716         { 24 } { X }
8717         { 25 } { Y }
8718         { 26 } { Z }
8719     }
8720 }

```

(End definition for `\int_to_alpha:n` and `\int_to_Alph:n`. These functions are documented on page 97.)

`\int_to_base:nn` Converting from base ten (#1) to a second base (#2) starts with computing #1: if it is
`\int_to_Base:nn` a complicated calculation, we shouldn't perform it twice. Then check the sign, store it,
`__int_to_base:nn` either - or `\c_empty_tl`, and feed the absolute value to the next auxiliary function.
`__int_to_Base:nn`
`__int_to_base:nnN` 8721 `\cs_new:Npn \int_to_base:nn #1`
`__int_to_Base:nnN` 8722 `{ \exp_args:Nf __int_to_base:nn { \int_eval:n {#1} } }`
`__int_to_base:nnnN` 8723 `\cs_new:Npn \int_to_Base:nn #1`
`__int_to_Base:nnnN` 8724 `{ \exp_args:Nf __int_to_Base:nn { \int_eval:n {#1} } }`
`__int_to_letter:n` 8725 `\cs_new:Npn __int_to_base:nn #1#2`
`__int_to_Letter:n` 8726 `{`
8727 `\int_compare:nNnTF {#1} < 0`
8728 `{ \exp_args:No __int_to_base:nnN { \use_none:n #1 } {#2} - }`
8729 `{ __int_to_base:nnN {#1} {#2} \c_empty_tl }`
8730 `}`
8731 `\cs_new:Npn __int_to_Base:nn #1#2`
8732 `{`
8733 `\int_compare:nNnTF {#1} < 0`
8734 `{ \exp_args:No __int_to_Base:nnN { \use_none:n #1 } {#2} - }`
8735 `{ __int_to_Base:nnN {#1} {#2} \c_empty_tl }`
8736 `}`

Here, the idea is to provide a recursive system to deal with the input. The output is built up after the end of the function. At each pass, the value in #1 is checked to see if it is less than the new base (#2). If it is, then it is converted directly, putting the sign back in front. On the other hand, if the value to convert is greater than or equal to the new base then the modulus and remainder values are found. The modulus is converted to a symbol and put on the right, and the remainder is carried forward to the next round.

8737 `\cs_new:Npn __int_to_base:nnN #1#2#3`
8738 `{`
8739 `\int_compare:nNnTF {#1} < {#2}`
8740 `{ \exp_last_unbraced:Nf #3 { __int_to_letter:n {#1} } }`
8741 `{`
8742 `\exp_args:Nf __int_to_base:nnnN`
8743 `{ __int_to_letter:n { \int_mod:nn {#1} {#2} } }`
8744 `{#1}`
8745 `{#2}`
8746 `#3`
8747 `}`
8748 `}`
8749 `\cs_new:Npn __int_to_base:nnnN #1#2#3#4`
8750 `{`
8751 `\exp_args:Nf __int_to_base:nnN`
8752 `{ \int_div_truncate:nn {#2} {#3} }`
8753 `{#3}`
8754 `#4`
8755 `#1`
8756 `}`
8757 `\cs_new:Npn __int_to_Base:nnN #1#2#3`
8758 `{`
8759 `\int_compare:nNnTF {#1} < {#2}`
8760 `{ \exp_last_unbraced:Nf #3 { __int_to_Letter:n {#1} } }`
8761 `{`
8762 `\exp_args:Nf __int_to_Base:nnnN`
8763 `{ __int_to_Letter:n { \int_mod:nn {#1} {#2} } }`
8764 `{#1}`

```

8765         {#2}
8766         #3
8767     }
8768 }
8769 \cs_new:Npn \__int_to_Base:nnnN #1#2#3#4
8770 {
8771     \exp_args:Nf \__int_to_Base:nnN
8772     { \int_div_truncate:nn {#2} {#3} }
8773     {#3}
8774     #4
8775     #1
8776 }

```

Convert to a letter only if necessary, otherwise simply return the value unchanged. It would be cleaner to use `\int_case:nn`, but in our case, the cases are contiguous, so it is forty times faster to use the `\if_case:w` primitive. The first `\exp_after:wN` expands the conditional, jumping to the correct case, the second one expands after the resulting character to close the conditional. Since `#1` might be an expression, and not directly a single digit, we need to evaluate it properly, and expand the trailing `\fi:`.

```

8777 \cs_new:Npn \__int_to_letter:n #1
8778 {
8779     \exp_after:wN \exp_after:wN
8780     \if_case:w \__int_eval:w #1 - 10 \__int_eval_end:
8781         a
8782     \or: b
8783     \or: c
8784     \or: d
8785     \or: e
8786     \or: f
8787     \or: g
8788     \or: h
8789     \or: i
8790     \or: j
8791     \or: k
8792     \or: l
8793     \or: m
8794     \or: n
8795     \or: o
8796     \or: p
8797     \or: q
8798     \or: r
8799     \or: s
8800     \or: t
8801     \or: u
8802     \or: v
8803     \or: w
8804     \or: x
8805     \or: y
8806     \or: z
8807     \else: \int_value:w \__int_eval:w #1 \exp_after:wN \__int_eval_end:
8808     \fi:
8809 }
8810 \cs_new:Npn \__int_to_Letter:n #1
8811 {

```

```

8812 \exp_after:wN \exp_after:wN
8813 \if_case:w \__int_eval:w #1 - 10 \__int_eval_end:
8814     A
8815 \or: B
8816 \or: C
8817 \or: D
8818 \or: E
8819 \or: F
8820 \or: G
8821 \or: H
8822 \or: I
8823 \or: J
8824 \or: K
8825 \or: L
8826 \or: M
8827 \or: N
8828 \or: O
8829 \or: P
8830 \or: Q
8831 \or: R
8832 \or: S
8833 \or: T
8834 \or: U
8835 \or: V
8836 \or: W
8837 \or: X
8838 \or: Y
8839 \or: Z
8840 \else: \int_value:w \__int_eval:w #1 \exp_after:wN \__int_eval_end:
8841 \fi:
8842 }

```

(End definition for `\int_to_base:nn` and others. These functions are documented on page 98.)

```

\int_to_bin:n  Wrappers around the generic function.
\int_to_hex:n  8843 \cs_new:Npn \int_to_bin:n #1
\int_to_Hex:n  8844 { \int_to_base:nn {#1} { 2 } }
\int_to_oct:n  8845 \cs_new:Npn \int_to_hex:n #1
                8846 { \int_to_base:nn {#1} { 16 } }
                8847 \cs_new:Npn \int_to_Hex:n #1
                8848 { \int_to_Base:nn {#1} { 16 } }
                8849 \cs_new:Npn \int_to_oct:n #1
                8850 { \int_to_base:nn {#1} { 8 } }

```

(End definition for `\int_to_bin:n` and others. These functions are documented on page 98.)

```

\int_to_roman:n The \__int_to_roman:w primitive creates tokens of category code 12 (other). Usually,
\int_to_Roman:n what is actually wanted is letters. The approach here is to convert the output of the
                  primitive into letters using appropriate control sequence names. That keeps everything
\__int_to_roman:N expandable. The loop is terminated by the conversion of the Q.
\__int_to_roman:N
\__int_to_roman_i:w 8851 \cs_new:Npn \int_to_roman:n #1
\__int_to_roman_v:w 8852 {
\__int_to_roman_x:w 8853 \exp_after:wN \__int_to_roman:N
\__int_to_roman_l:w 8854 \__int_to_roman:w \int_eval:n {#1} Q
\__int_to_roman_c:w
\__int_to_roman_d:w
\__int_to_roman_m:w
\__int_to_roman_Q:w
\__int_to_Roman_i:w
\__int_to_Roman_v:w
\__int_to_Roman_x:w
\__int_to_Roman_l:w
\__int_to_Roman_c:w
\__int_to_Roman_d:w

```

```

8855 }
8856 \cs_new:Npn \__int_to_roman:N #1
8857 {
8858   \use:c { __int_to_roman_ #1 :w }
8859   \__int_to_roman:N
8860 }
8861 \cs_new:Npn \int_to_Roman:n #1
8862 {
8863   \exp_after:wN \__int_to_Roman_aux:N
8864   \__int_to_roman:w \int_eval:n {#1} Q
8865 }
8866 \cs_new:Npn \__int_to_Roman_aux:N #1
8867 {
8868   \use:c { __int_to_Roman_ #1 :w }
8869   \__int_to_Roman_aux:N
8870 }
8871 \cs_new:Npn \__int_to_roman_i:w { i }
8872 \cs_new:Npn \__int_to_roman_v:w { v }
8873 \cs_new:Npn \__int_to_roman_x:w { x }
8874 \cs_new:Npn \__int_to_roman_l:w { l }
8875 \cs_new:Npn \__int_to_roman_c:w { c }
8876 \cs_new:Npn \__int_to_roman_d:w { d }
8877 \cs_new:Npn \__int_to_roman_m:w { m }
8878 \cs_new:Npn \__int_to_roman_Q:w #1 { }
8879 \cs_new:Npn \__int_to_Roman_i:w { I }
8880 \cs_new:Npn \__int_to_Roman_v:w { V }
8881 \cs_new:Npn \__int_to_Roman_x:w { X }
8882 \cs_new:Npn \__int_to_Roman_l:w { L }
8883 \cs_new:Npn \__int_to_Roman_c:w { C }
8884 \cs_new:Npn \__int_to_Roman_d:w { D }
8885 \cs_new:Npn \__int_to_Roman_m:w { M }
8886 \cs_new:Npn \__int_to_Roman_Q:w #1 { }

```

(End definition for `\int_to_roman:n` and others. These functions are documented on page 98.)

12.9 Converting from other formats to integers

`__int_pass_signs:wn` `__int_pass_signs_end:wn` Called as `__int_pass_signs:wn <signs and digits> \s__int_stop {<code>}`, this function leaves in the input stream any sign it finds, then inserts the `<code>` before the first non-sign token (and removes `\s__int_stop`). More precisely, it deletes any + and passes any - to the input stream, hence should be called in an integer expression.

```

8887 \cs_new:Npn \__int_pass_signs:wn #1
8888 {
8889   \if:w + \if:w - \exp_not:N #1 + \fi: \exp_not:N #1
8890   \exp_after:wN \__int_pass_signs:wn
8891   \else:
8892     \exp_after:wN \__int_pass_signs_end:wn
8893     \exp_after:wN #1
8894   \fi:
8895 }
8896 \cs_new:Npn \__int_pass_signs_end:wn #1 \s__int_stop #2 { #2 #1 }

```

(End definition for `__int_pass_signs:wn` and `__int_pass_signs_end:wn`.)

`\int_from_alph:n` First take care of signs then loop through the input using the `recursion` quarks. The `__int_from_alph:nN` auxiliary collects in its first argument the value obtained so far, and the auxiliary `__int_from_alph:N` converts one letter to an expression which evaluates to the correct number.

```

8897 \cs_new:Npn \int_from_alph:n #1
8898 {
8899   \int_eval:n
8900   {
8901     \exp_after:wN \__int_pass_signs:wn \tl_to_str:n {#1}
8902     \s__int_stop { \__int_from_alph:nN { 0 } }
8903     \q__int_recursion_tail \q__int_recursion_stop
8904   }
8905 }
8906 \cs_new:Npn \__int_from_alph:nN #1#2
8907 {
8908   \__int_if_recursion_tail_stop_do:Nn #2 {#1}
8909   \exp_args:Nf \__int_from_alph:nN
8910   { \int_eval:n { #1 * 26 + \__int_from_alph:N #2 } }
8911 }
8912 \cs_new:Npn \__int_from_alph:N #1
8913 { '#1 - \int_compare:nNnTF { '#1 } < { 91 } { 64 } { 96 } }

```

(End definition for `\int_from_alph:n`, `__int_from_alph:nN`, and `__int_from_alph:N`. This function is documented on page 98.)

`\int_from_base:nn` Leave the signs into the integer expression, then loop through characters, collecting the value found so far in the first argument of `__int_from_base:nnN`. To convert a single character, `__int_from_base:N` checks first for digits, then distinguishes lower from upper case letters, turning them into the appropriate number. Note that this auxiliary does not use `\int_eval:n`, hence is not safe for general use.

```

8914 \cs_new:Npn \int_from_base:nn #1#2
8915 {
8916   \int_eval:n
8917   {
8918     \exp_after:wN \__int_pass_signs:wn \tl_to_str:n {#1}
8919     \s__int_stop { \__int_from_base:nnN { 0 } {#2} }
8920     \q__int_recursion_tail \q__int_recursion_stop
8921   }
8922 }
8923 \cs_new:Npn \__int_from_base:nnN #1#2#3
8924 {
8925   \__int_if_recursion_tail_stop_do:Nn #3 {#1}
8926   \exp_args:Nf \__int_from_base:nnN
8927   { \int_eval:n { #1 * #2 + \__int_from_base:N #3 } }
8928   {#2}
8929 }
8930 \cs_new:Npn \__int_from_base:N #1
8931 {
8932   \int_compare:nNnTF { '#1 } < { 58 }
8933   {#1}
8934   { '#1 - \int_compare:nNnTF { '#1 } < { 91 } { 55 } { 87 } }
8935 }

```

(End definition for `\int_from_base:nn`, `__int_from_base:nnN`, and `__int_from_base:N`. This function is documented on page 99.)

`\int_from_bin:n` Wrappers around the generic function.

```

\int_from_hex:n
\int_from_oct:n

8936 \cs_new:Npn \int_from_bin:n #1
8937 { \int_from_base:nn {#1} { 2 } }
8938 \cs_new:Npn \int_from_hex:n #1
8939 { \int_from_base:nn {#1} { 16 } }
8940 \cs_new:Npn \int_from_oct:n #1
8941 { \int_from_base:nn {#1} { 8 } }

```

(End definition for `\int_from_bin:n`, `\int_from_hex:n`, and `\int_from_oct:n`. These functions are documented on page 99.)

`\c_int_from_roman_i_int` Constants used to convert from Roman numerals to integers.

```

\c_int_from_roman_v_int
\c_int_from_roman_x_int
\c_int_from_roman_l_int
\c_int_from_roman_c_int
\c_int_from_roman_d_int
\c_int_from_roman_m_int
\c_int_from_roman_I_int
\c_int_from_roman_V_int
\c_int_from_roman_X_int
\c_int_from_roman_L_int
\c_int_from_roman_C_int
\c_int_from_roman_D_int
\c_int_from_roman_M_int

8942 \int_const:cn { c__int_from_roman_i_int } { 1 }
8943 \int_const:cn { c__int_from_roman_v_int } { 5 }
8944 \int_const:cn { c__int_from_roman_x_int } { 10 }
8945 \int_const:cn { c__int_from_roman_l_int } { 50 }
8946 \int_const:cn { c__int_from_roman_c_int } { 100 }
8947 \int_const:cn { c__int_from_roman_d_int } { 500 }
8948 \int_const:cn { c__int_from_roman_m_int } { 1000 }
8949 \int_const:cn { c__int_from_roman_I_int } { 1 }
8950 \int_const:cn { c__int_from_roman_V_int } { 5 }
8951 \int_const:cn { c__int_from_roman_X_int } { 10 }
8952 \int_const:cn { c__int_from_roman_L_int } { 50 }
8953 \int_const:cn { c__int_from_roman_C_int } { 100 }
8954 \int_const:cn { c__int_from_roman_D_int } { 500 }
8955 \int_const:cn { c__int_from_roman_M_int } { 1000 }

```

(End definition for `\c_int_from_roman_i_int` and others.)

`\int_from_roman:n` The method here is to iterate through the input, finding the appropriate value for each letter and building up a sum. This is then evaluated by \TeX . If any unknown letter is found, skip to the closing parenthesis and insert `*0-1` afterwards, to replace the value by -1 .

```

__int_from_roman:NN
__int_from_roman_error:w

8956 \cs_new:Npn \int_from_roman:n #1
8957 {
8958   \int_eval:n
8959   {
8960     (
8961       0
8962       \exp_after:wN __int_from_roman:NN \tl_to_str:n {#1}
8963       \q__int_recursion_tail \q__int_recursion_tail \q__int_recursion_stop
8964     )
8965   }
8966 }
8967 \cs_new:Npn __int_from_roman:NN #1#2
8968 {
8969   \__int_if_recursion_tail_stop:N #1
8970   \int_if_exist:cF { c__int_from_roman_ #1 _int }
8971   { \__int_from_roman_error:w }
8972   \__int_if_recursion_tail_stop_do:Nn #2
8973   { + \use:c { c__int_from_roman_ #1 _int } }
8974   \int_if_exist:cF { c__int_from_roman_ #2 _int }
8975   { \__int_from_roman_error:w }
8976   \int_compare:nNnTF

```

```

8977 { \use:c { c__int_from_roman_ #1 _int } }
8978 <
8979 { \use:c { c__int_from_roman_ #2 _int } }
8980 {
8981 + \use:c { c__int_from_roman_ #2 _int }
8982 - \use:c { c__int_from_roman_ #1 _int }
8983 \__int_from_roman:NN
8984 }
8985 {
8986 + \use:c { c__int_from_roman_ #1 _int }
8987 \__int_from_roman:NN #2
8988 }
8989 }
8990 \cs_new:Npn \__int_from_roman_error:w #1 \q__int_recursion_stop #2
8991 { #2 * 0 - 1 }

```

(End definition for `\int_from_roman:n`, `__int_from_roman:NN`, and `__int_from_roman_error:w`. This function is documented on page 99.)

12.10 Viewing integer

`\int_show:N` Diagnostics.
`\int_show:c` 8992 \cs_new_eq:NN \int_show:N __kernel_register_show:N
`__int_show:nN` 8993 \cs_generate_variant:Nn \int_show:N { c }

(End definition for `\int_show:N` and `__int_show:nN`. This function is documented on page 100.)

`\int_show:n` We don't use the TeX primitive `\showthe` to show integer expressions: this gives a more unified output.

```

8994 \cs_new_protected:Npn \int_show:n
8995 { \msg_show_eval:Nn \int_eval:n }

```

(End definition for `\int_show:n`. This function is documented on page 100.)

`\int_log:N` Diagnostics.
`\int_log:c` 8996 \cs_new_eq:NN \int_log:N __kernel_register_log:N
8997 \cs_generate_variant:Nn \int_log:N { c }

(End definition for `\int_log:N`. This function is documented on page 100.)

`\int_log:n` Similar to `\int_show:n`.

```

8998 \cs_new_protected:Npn \int_log:n
8999 { \msg_log_eval:Nn \int_eval:n }

```

(End definition for `\int_log:n`. This function is documented on page 100.)

12.11 Random integers

`\int_rand:nn` Defined in `l3fp-random`.

(End definition for `\int_rand:nn`. This function is documented on page 99.)

12.12 Constant integers

`\c_zero_int` The zero is defined in `l3basics`.

`\c_one_int` 9000 `\int_const:Nn \c_one_int { 1 }`

(End definition for `\c_zero_int` and `\c_one_int`. These variables are documented on page 100.)

`\c_max_int` The largest number allowed is $2^{31} - 1$

9001 `\int_const:Nn \c_max_int { 2 147 483 647 }`

(End definition for `\c_max_int`. This variable is documented on page 100.)

`\c_max_char_int` The largest character code is 1114111 (hexadecimal 10FFFF) in XeTeX and LuaTeX and 255 in other engines. In many places pTeX and upTeX support larger character codes but for instance the values of `\lccode` are restricted to $[0, 255]$.

```
9002 \int_const:Nn \c_max_char_int
9003 {
9004   \if_int_odd:w 0
9005     \cs_if_exist:NT \tex luatexversion:D { 1 }
9006     \cs_if_exist:NT \tex XeTeXversion:D { 1 } ~
9007     "10FFFF
9008   \else:
9009     "FF
9010   \fi:
9011 }
```

(End definition for `\c_max_char_int`. This variable is documented on page 100.)

12.13 Scratch integers

`\l_tmpa_int` We provide two local and two global scratch counters, maybe we need more or less.

`\l_tmpb_int` 9012 `\int_new:N \l_tmpa_int`

`\g_tmpa_int` 9013 `\int_new:N \l_tmpb_int`

`\g_tmpb_int` 9014 `\int_new:N \g_tmpa_int`

9015 `\int_new:N \g_tmpb_int`

(End definition for `\l_tmpa_int` and others. These variables are documented on page 100.)

12.14 Integers for earlier modules

<@@=seq>

`\l__int_internal_a_int`

`\l__int_internal_b_int`

9016 `\int_new:N \l__int_internal_a_int`

9017 `\int_new:N \l__int_internal_b_int`

(End definition for `\l__int_internal_a_int` and `\l__int_internal_b_int`.)

9018 `</package>`

13 l3flag implementation

9019 `\package`

9020 `\@@=flag`

The following test files are used for this code: m3flag001.

13.1 Non-expandable flag commands

The height h of a flag (initially zero) is stored by setting control sequences of the form `\flag <name> <integer>` to `\relax` for $0 \leq \langle integer \rangle < h$. When a flag is raised, a “trap” function `\flag <name>` is called. The existence of this function is also used to test for the existence of a flag.

`\flag_new:n` For each flag, we define a “trap” function, which by default simply increases the flag by 1 by letting the appropriate control sequence to `\relax`. This can be done expandably!

```
9021 \cs_new_protected:Npn \flag_new:n #1
9022 {
9023   \cs_new:cpn { flag~#1 } ##1 ;
9024   { \exp_after:wN \use_none:n \cs:w flag~#1~##1 \cs_end: }
9025 }
```

(End definition for \flag_new:n. This function is documented on page 103.)

`\flag_clear:n` `__flag_clear:wn` Undefine control sequences, starting from the 0 flag, upwards, until reaching an undefined control sequence. We don’t use `\cs_undefine:c` because that would act globally. When the option `check-declarations` is used, check for the function defined by `\flag_new:n`.

```
9026 \cs_new_protected:Npn \flag_clear:n #1 { \__flag_clear:wn 0 ; {#1} }
9027 \cs_new_protected:Npn \__flag_clear:wn #1 ; #2
9028 {
9029   \if_cs_exist:w flag~#2~#1 \cs_end:
9030   \cs_set_eq:cN { flag~#2~#1 } \tex_undefined:D
9031   \exp_after:wN \__flag_clear:wn
9032   \int_value:w \int_eval:w 1 + #1
9033   \else:
9034     \use_i:nnn
9035   \fi:
9036   ; {#2}
9037 }
```

(End definition for \flag_clear:n and __flag_clear:wn. This function is documented on page 103.)

`\flag_clear_new:n` As for other datatypes, clear the *<flag>* or create a new one, as appropriate.

```
9038 \cs_new_protected:Npn \flag_clear_new:n #1
9039 { \flag_if_exist:nTF {#1} { \flag_clear:n } { \flag_new:n } {#1} }
```

(End definition for \flag_clear_new:n. This function is documented on page 103.)

`\flag_show:n` `\flag_log:n` `__flag_show:Nn` Show the height (terminal or log file) using appropriate `l3msg` auxiliaries.

```
9040 \cs_new_protected:Npn \flag_show:n { \__flag_show:Nn \tl_show:n }
9041 \cs_new_protected:Npn \flag_log:n { \__flag_show:Nn \tl_log:n }
9042 \cs_new_protected:Npn \__flag_show:Nn #1#2
9043 {
9044   \exp_args:Nc \__kernel_chk_defined:NT { flag~#2 }
9045   {
```

```

9046         \exp_args:Nx #1
9047         { \tl_to_str:n { flag~#2~height } = \flag_height:n {#2} }
9048     }
9049 }

```

(End definition for `\flag_show:n`, `\flag_log:n`, and `__flag_show:Nn`. These functions are documented on page 103.)

13.2 Expandable flag commands

`\flag_if_exist_p:n` A flag exist if the corresponding trap `\flag <flag name>:n` is defined.

```

\flag_if_exist:nTF
9050 \prg_new_conditional:Npnn \flag_if_exist:n #1 { p , T , F , TF }
9051 {
9052     \cs_if_exist:cTF { flag~#1 }
9053     { \prg_return_true: } { \prg_return_false: }
9054 }

```

(End definition for `\flag_if_exist:nTF`. This function is documented on page 104.)

`\flag_if_raised_p:n` Test if the flag has a non-zero height, by checking the 0 control sequence.

```

\flag_if_raised:nTF
9055 \prg_new_conditional:Npnn \flag_if_raised:n #1 { p , T , F , TF }
9056 {
9057     \if_cs_exist:w flag~#1~0 \cs_end:
9058     \prg_return_true:
9059     \else:
9060     \prg_return_false:
9061     \fi:
9062 }

```

(End definition for `\flag_if_raised:nTF`. This function is documented on page 104.)

`\flag_height:n` Extract the value of the flag by going through all of the control sequences starting from 0.

```

\__flag_height_loop:wn
\__flag_height_end:wn
9063 \cs_new:Npn \flag_height:n #1 { \__flag_height_loop:wn 0; {#1} }
9064 \cs_new:Npn \__flag_height_loop:wn #1 ; #2
9065 {
9066     \if_cs_exist:w flag~#2~#1 \cs_end:
9067     \exp_after:wN \__flag_height_loop:wn \int_value:w \int_eval:w 1 +
9068     \else:
9069     \exp_after:wN \__flag_height_end:wn
9070     \fi:
9071     #1 ; {#2}
9072 }
9073 \cs_new:Npn \__flag_height_end:wn #1 ; #2 {#1}

```

(End definition for `\flag_height:n`, `__flag_height_loop:wn`, and `__flag_height_end:wn`. This function is documented on page 104.)

`\flag_raise:n` Simply apply the trap to the height, after expanding the latter.

```

9074 \cs_new:Npn \flag_raise:n #1
9075 {
9076     \cs:w flag~#1 \exp_after:wN \cs_end:
9077     \int_value:w \flag_height:n {#1} ;
9078 }

```

(End definition for `\flag_raise:n`. This function is documented on page 104.)

```

9079 </package>

```

14 l3prg implementation

The following test files are used for this code: *m3prg001.lvt,m3prg002.lvt,m3prg003.lvt.*

```
9080 <*package>
```

14.1 Primitive conditionals

`\if_bool:N` Those two primitive T_EX conditionals are synonyms. `\if_bool:N` is defined in l3basics, as it's needed earlier to define quark test functions.

```
9081 \cs_new_eq:NN \if_predicate:w \tex_ifodd:D
```

(End definition for `\if_bool:N` and `\if_predicate:w`. These functions are documented on page 113.)

14.2 Defining a set of conditional functions

These are all defined in l3basics, as they are needed “early”. This is just a reminder!

(End definition for `\prg_set_conditional:Npnn` and others. These functions are documented on page 105.)

14.3 The boolean data type

```
9082 <@@=bool>
```

Boolean variables have to be initiated when they are created. Other than that there is not much to say here.

```
9083 \cs_new_protected:Npn \bool_new:N #1 { \cs_new_eq:NN #1 \c_false_bool }
9084 \cs_generate_variant:Nn \bool_new:N { c }
```

(End definition for `\bool_new:N`. This function is documented on page 107.)

`\bool_const:Nn` A merger between `\tl_const:Nn` and `\bool_set:Nn`.

```
\bool_const:cn
9085 \cs_new_protected:Npn \bool_const:Nn #1#2
9086 {
9087   \__kernel_chk_if_free_cs:N #1
9088   \tex_global:D \tex_chardef:D #1 = \bool_if_p:n {#2}
9089 }
9090 \cs_generate_variant:Nn \bool_const:Nn { c }
```

(End definition for `\bool_const:Nn`. This function is documented on page 108.)

`\bool_set_true:N` Setting is already pretty easy. When `check-declarations` is active, the definitions are patched to make sure the boolean exists. This is needed because booleans are not based on token lists nor on T_EX registers.

```
\bool_set_true:c
\bool_gset_true:N
\bool_gset_true:c
\bool_set_false:N
\bool_set_false:c
\bool_gset_false:N
\bool_gset_false:c
9091 \cs_new_protected:Npn \bool_set_true:N #1
9092 { \cs_set_eq:NN #1 \c_true_bool }
9093 \cs_new_protected:Npn \bool_set_false:N #1
9094 { \cs_set_eq:NN #1 \c_false_bool }
9095 \cs_new_protected:Npn \bool_gset_true:N #1
9096 { \cs_gset_eq:NN #1 \c_true_bool }
9097 \cs_new_protected:Npn \bool_gset_false:N #1
9098 { \cs_gset_eq:NN #1 \c_false_bool }
9099 \cs_generate_variant:Nn \bool_set_true:N { c }
9100 \cs_generate_variant:Nn \bool_set_false:N { c }
9101 \cs_generate_variant:Nn \bool_gset_true:N { c }
9102 \cs_generate_variant:Nn \bool_gset_false:N { c }
```

(End definition for `\bool_set_true:N` and others. These functions are documented on page 108.)

`\bool_set_eq:NN` The usual copy code. While it would be cleaner semantically to copy the `\cs_set_eq:NN` family of functions, we copy `\tl_set_eq:NN` because that has the correct checking code.

`\bool_set_eq:cN`

`\bool_set_eq:Nc` 9103 `\cs_new_eq:NN \bool_set_eq:NN \tl_set_eq:NN`

`\bool_set_eq:cc` 9104 `\cs_new_eq:NN \bool_gset_eq:NN \tl_gset_eq:NN`

`\bool_gset_eq:NN` 9105 `\cs_generate_variant:Nn \bool_set_eq:NN { Nc, cN, cc }`

`\bool_gset_eq:cN` 9106 `\cs_generate_variant:Nn \bool_gset_eq:NN { Nc, cN, cc }`

`\bool_gset_eq:Nc` (End definition for `\bool_set_eq:NN` and `\bool_gset_eq:NN`. These functions are documented on page 108.)

`\bool_gset_eq:cc`

`\bool_set:Nn` This function evaluates a boolean expression and assigns the first argument the meaning

`\bool_set:cn` `\c_true_bool` or `\c_false_bool`. Again, we include some checking code. It is important

`\bool_gset:Nn` to evaluate the expression before applying the `\chardef` primitive, because that primitive

`\bool_gset:cn` sets the left-hand side to `\scan_stop:` before looking for the right-hand side.

9107 `\cs_new_protected:Npn \bool_set:Nn #1#2`

9108 `{`

9109 `\exp_last_unbraced:NNNf`

9110 `\tex_chardef:D #1 = { \bool_if_p:n {#2} }`

9111 `}`

9112 `\cs_new_protected:Npn \bool_gset:Nn #1#2`

9113 `{`

9114 `\exp_last_unbraced:NNNNf`

9115 `\tex_global:D \tex_chardef:D #1 = { \bool_if_p:n {#2} }`

9116 `}`

9117 `\cs_generate_variant:Nn \bool_set:Nn { c }`

9118 `\cs_generate_variant:Nn \bool_gset:Nn { c }`

(End definition for `\bool_set:Nn` and `\bool_gset:Nn`. These functions are documented on page 108.)

14.4 Internal auxiliaries

`\q__bool_recursion_tail` Internal recursion quarks.

`\q__bool_recursion_stop` 9119 `\quark_new:N \q__bool_recursion_tail`

9120 `\quark_new:N \q__bool_recursion_stop`

(End definition for `\q__bool_recursion_tail` and `\q__bool_recursion_stop`.)

`__bool_use_i_delimit_by_q_recursion_stop:nw` Functions to gobble up to a quark.

9121 `\cs_new:Npn __bool_use_i_delimit_by_q_recursion_stop:nw`

9122 `#1 #2 \q__bool_recursion_stop {#1}`

(End definition for `__bool_use_i_delimit_by_q_recursion_stop:nw`.)

`__bool_if_recursion_tail_stop_do:nn` Functions to query recursion quarks.

9123 `__kernel_quark_new_test:N __bool_if_recursion_tail_stop_do:nn`

(End definition for `__bool_if_recursion_tail_stop_do:nn`.)

`\bool_if_p:N` Straight forward here. We could optimize here if we wanted to as the boolean can just be input directly.

```

\bool_if_p:c
\bool_if:N $\underline{TF}$ 
\bool_if:c $\underline{TF}$ 
9124 \prg_new_conditional:Npnn \bool_if:N #1 { p , T , F , TF }
9125 {
9126   \if_bool:N #1
9127   \prg_return_true:
9128   \else:
9129   \prg_return_false:
9130   \fi:
9131 }
9132 \prg_generate_conditional_variant:Nnn \bool_if:N { c } { p , T , F , TF }
```

(End definition for `\bool_if:N \underline{TF}` . This function is documented on page 108.)

`\bool_show:n` Show the truth value of the boolean, as true or false.

```

\bool_log:n
\__bool_to_str:n
9133 \cs_new_protected:Npn \bool_show:n
9134 { \msg_show_eval:Nn \__bool_to_str:n }
9135 \cs_new_protected:Npn \bool_log:n
9136 { \msg_log_eval:Nn \__bool_to_str:n }
9137 \cs_new:Npn \__bool_to_str:n #1
9138 { \bool_if:nTF {#1} { true } { false } }
```

(End definition for `\bool_show:n`, `\bool_log:n`, and `__bool_to_str:n`. These functions are documented on page 108.)

`\bool_show:N` Show the truth value of the boolean, as true or false.

```

\bool_show:c
\bool_log:N
\bool_log:c
\__bool_show:NN
9139 \cs_new_protected:Npn \bool_show:N { \__bool_show:NN \tl_show:n }
9140 \cs_generate_variant:Nn \bool_show:N { c }
9141 \cs_new_protected:Npn \bool_log:N { \__bool_show:NN \tl_log:n }
9142 \cs_generate_variant:Nn \bool_log:N { c }
9143 \cs_new_protected:Npn \__bool_show:NN #1#2
9144 {
9145   \__kernel_chk_defined:NT #2
9146   { \exp_args:Nx #1 { \token_to_str:N #2 = \__bool_to_str:n {#2} } }
9147 }
```

(End definition for `\bool_show:N`, `\bool_log:N`, and `__bool_show:NN`. These functions are documented on page 108.)

`\l_tmpa_bool` A few booleans just if you need them.

```

\l_tmpb_bool
\g_tmpa_bool
\g_tmpb_bool
9148 \bool_new:N \l_tmpa_bool
9149 \bool_new:N \l_tmpb_bool
9150 \bool_new:N \g_tmpa_bool
9151 \bool_new:N \g_tmpb_bool
```

(End definition for `\l_tmpa_bool` and others. These variables are documented on page 109.)

`\bool_if_exist_p:N` Copies of the cs functions defined in l3basics.

```

\bool_if_exist_p:c
\bool_if_exist:N $\underline{TF}$ 
\bool_if_exist:c $\underline{TF}$ 
9152 \prg_new_eq_conditional:NNn \bool_if_exist:N \cs_if_exist:N
9153 { TF , T , F , p }
9154 \prg_new_eq_conditional:NNn \bool_if_exist:c \cs_if_exist:c
9155 { TF , T , F , p }
```

(End definition for `\bool_if_exist:N \underline{TF}` . This function is documented on page 109.)

14.5 Boolean expressions

`\bool_if_p:n` Evaluating the truth value of a list of predicates is done using an input syntax somewhat similar to the one found in other programming languages with (and) for grouping, ! for logical “Not”, && for logical “And” and || for logical “Or”. However, they perform eager evaluation. We shall use the terms Not, And, Or, Open and Close for these operations.

`\bool_if:nTF`

Any expression is terminated by a Close operation. Evaluation happens from left to right in the following manner using a GetNext function:

- If an Open is seen, start evaluating a new expression using the Eval function and call GetNext again.
- If a Not is seen, remove the ! and call a GetNext function with the logic reversed.
- If none of the above, reinsert the token found (this is supposed to be a predicate function) in front of an Eval function, which evaluates it to the boolean value $\langle true \rangle$ or $\langle false \rangle$.

The Eval function then contains a post-processing operation which grabs the instruction following the predicate. This is either And, Or or Close. In each case the truth value is used to determine where to go next. The following situations can arise:

$\langle true \rangle$ **And** Current truth value is true, logical And seen, continue with GetNext to examine truth value of next boolean (sub-)expression.

$\langle false \rangle$ **And** Current truth value is false, logical And seen, stop using the values of predicates within this sub-expression until the next Close. Then return $\langle false \rangle$.

$\langle true \rangle$ **Or** Current truth value is true, logical Or seen, stop using the values of predicates within this sub-expression until the nearest Close. Then return $\langle true \rangle$.

$\langle false \rangle$ **Or** Current truth value is false, logical Or seen, continue with GetNext to examine truth value of next boolean (sub-)expression.

$\langle true \rangle$ **Close** Current truth value is true, Close seen, return $\langle true \rangle$.

$\langle false \rangle$ **Close** Current truth value is false, Close seen, return $\langle false \rangle$.

```

9156 \prg_new_conditional:Npnn \bool_if:n #1 { T , F , TF }
9157 {
9158   \if_predicate:w \bool_if_p:n {#1}
9159   \prg_return_true:
9160   \else:
9161   \prg_return_false:
9162   \fi:
9163 }
```

(End definition for `\bool_if:nTF`. This function is documented on page 110.)

`\bool_if_p:n` To speed up the case of a single predicate, f-expand and check whether the result is one token (possibly surrounded by spaces), which must be `\c_true_bool` or `\c_false_bool`. We use a version of `\tl_if_single:nTF` optimized for speed since we know that an empty #1 is an error. The auxiliary `__bool_if_p_aux:w` removes the trailing parenthesis and gets rid of any space. For the general case, first issue a `\group_align_safe_begin:` as we are using && as syntax shorthand for the And operation and we need to hide it for

`__bool_if_p:n`

`__bool_if_p_aux:w`

TeX. This group is closed after `__bool_get_next:NN` returns `\c_true_bool` or `\c_false_bool`. That function requires the trailing parenthesis to know where the expression ends.

```

9164 \cs_new:Npn \bool_if_p:n { \exp_args:Nf \__bool_if_p:n }
9165 \cs_new:Npn \__bool_if_p:n #1
9166 {
9167   \tl_if_empty:oT { \use_none:nn #1 . } { \__bool_if_p_aux:w }
9168   \group_align_safe_begin:
9169   \exp_after:wN
9170   \group_align_safe_end:
9171   \exp:w \exp_end_continue_f:w % (
9172   \__bool_get_next:NN \use_i:nnnn #1 )
9173 }
9174 \cs_new:Npn \__bool_if_p_aux:w #1 \use_i:nnnn #2#3 {#2}

```

(End definition for `\bool_if_p:n`, `__bool_if_p:n`, and `__bool_if_p_aux:w`. This function is documented on page 110.)

`__bool_get_next:NN` The GetNext operation. Its first argument is `\use_i:nnnn`, `\use_ii:nnnn`, `\use_iii:nnnn`, or `\use_iv:nnnn` (we call these “states”). In the first state, this function eventually expand to the truth value `\c_true_bool` or `\c_false_bool` of the expression which follows until the next unmatched closing parenthesis. For instance “`__bool_get_next:NN \use_i:nnnn \c_true_bool && \c_true_bool)`” (including the closing parenthesis) expands to `\c_true_bool`. In the second state (after a `!`) the logic is reversed. We call these two states “normal” and the next two “skipping”. In the third state (after `\c_true_bool||`) it always returns `\c_true_bool`. In the fourth state (after `\c_false_bool&&`) it always returns `\c_false_bool` and also stops when encountering `||`, not only parentheses. This code itself is a switch: if what follows is neither `!` nor `(`, we assume it is a predicate.

```

9175 \cs_new:Npn \__bool_get_next:NN #1#2
9176 {
9177   \use:c
9178   {
9179     __bool_
9180     \if_meaning:w !#2 ! \else: \if_meaning:w (#2 ( \else: p \fi: \fi:
9181     :Nw
9182   }
9183   #1 #2
9184 }

```

(End definition for `__bool_get_next:NN`.)

`__bool_!:Nw` The Not operation reverses the logic: it discards the `!` token and calls the GetNext operation with the appropriate first argument. Namely the first and second states are interchanged, but after `\c_true_bool||` or `\c_false_bool&&` the `!` is ignored.

```

9185 \cs_new:cpn { __bool_!:Nw } #1#2
9186 {
9187   \exp_after:wN \__bool_get_next:NN
9188   #1 \use_ii:nnnn \use_i:nnnn \use_iii:nnnn \use_iv:nnnn
9189 }

```

(End definition for `__bool_!:Nw`.)

`__bool_(:Nw` The Open operation starts a sub-expression after discarding the open parenthesis. This is done by calling `GetNext` (which eventually discards the corresponding closing parenthesis), with a post-processing step which looks for And, Or or Close after the group.

```

9190 \cs_new:cpn { __bool_(:Nw } #1#2
9191 {
9192   \exp_after:wN \__bool_choose:NNN \exp_after:wN #1
9193   \int_value:w \__bool_get_next:NN \use_i:nnnn
9194 }

```

(End definition for `__bool_(:Nw`.)

`__bool_p:Nw` If what follows `GetNext` is neither `!` nor `(`, evaluate the predicate using the primitive `\int_value:w`. The canonical `true` and `false` values have numerical values 1 and 0 respectively. Look for And, Or or Close afterwards.

```

9195 \cs_new:cpn { __bool_p:Nw } #1
9196 { \exp_after:wN \__bool_choose:NNN \exp_after:wN #1 \int_value:w }

```

(End definition for `__bool_p:Nw`.)

`__bool_choose:NNN` The arguments are `#1`: a function such as `\use_i:nnnn`, `#2`: 0 or 1 encoding the current truth value, `#3`: the next operation, And, Or or Close. We distinguish three cases according to a combination of `#1` and `#2`. Case 2 is when `#1` is `\use_iii:nnnn` (state 3), namely after `\c_true_bool ||`. Case 1 is when `#1` is `\use_i:nnnn` and `#2` is `true` or when `#1` is `\use_ii:nnnn` and `#2` is `false`, for instance for `!\c_false_bool`. Case 0 includes the same with `true/false` interchanged and the case where `#1` is `\use_iv:nnnn` namely after `\c_false_bool &&`.

`__bool_|_0:` When seeing `)` the current subexpression is done, leave the appropriate boolean.
`__bool_|_1:` When seeing `&` in case 0 go into state 4, equivalent to having seen `\c_false_bool &&`.
`__bool_|_2:` In case 1, namely when the argument is `true` and we are in a normal state continue in the normal state 1. In case 2, namely when skipping alternatives in an Or, continue in the same state. When seeing `|` in case 0, continue in a normal state; in particular stop skipping for `\c_false_bool &&` because that binds more tightly than `||`. In the other two cases start skipping for `\c_true_bool ||`.

```

9197 \cs_new:Npn \__bool_choose:NNN #1#2#3
9198 {
9199   \use:c
9200   {
9201     __bool_ \token_to_str:N #3 _
9202     #1 #2 { \if_meaning:w 0 #2 1 \else: 0 \fi: } 2 0 :
9203   }
9204 }
9205 \cs_new:cpn { __bool_)_0: } { \c_false_bool }
9206 \cs_new:cpn { __bool_)_1: } { \c_true_bool }
9207 \cs_new:cpn { __bool_)_2: } { \c_true_bool }
9208 \cs_new:cpn { __bool_&_0: } & { \__bool_get_next:NN \use_iv:nnnn }
9209 \cs_new:cpn { __bool_&_1: } & { \__bool_get_next:NN \use_i:nnnn }
9210 \cs_new:cpn { __bool_&_2: } & { \__bool_get_next:NN \use_iii:nnnn }
9211 \cs_new:cpn { __bool_|_0: } | { \__bool_get_next:NN \use_i:nnnn }
9212 \cs_new:cpn { __bool_|_1: } | { \__bool_get_next:NN \use_iii:nnnn }
9213 \cs_new:cpn { __bool_|_2: } | { \__bool_get_next:NN \use_iii:nnnn }

```

(End definition for `__bool_choose:NNN` and others.)

\bool_lazy_all_p:n Go through the list of expressions, stopping whenever an expression is **false**. If the end is reached without finding any **false** expression, then the result is **true**.

\bool_lazy_all:nTF

```

9214 \cs_new:Npn \bool_lazy_all_p:n #1
9215 { \__bool_lazy_all:n #1 \q__bool_recursion_tail \q__bool_recursion_stop }
9216 \prg_new_conditional:Npnn \bool_lazy_all:n #1 { T , F , TF }
9217 {
9218   \if_predicate:w \bool_lazy_all_p:n {#1}
9219   \prg_return_true:
9220   \else:
9221   \prg_return_false:
9222   \fi:
9223 }
9224 \cs_new:Npn \__bool_lazy_all:n #1
9225 {
9226   \__bool_if_recursion_tail_stop_do:nn {#1} { \c_true_bool }
9227   \bool_if:nF {#1}
9228   { \__bool_use_i_delimit_by_q_recursion_stop:nw { \c_false_bool } }
9229   \__bool_lazy_all:n
9230 }

```

(End definition for **\bool_lazy_all:nTF** and **__bool_lazy_all:n**. This function is documented on page 110.)

\bool_lazy_and_p:nn Only evaluate the second expression if the first is **true**. Note that #2 must be removed as an argument, not just by skipping to the **\else:** branch of the conditional since #2 may contain unbalanced **TeX** conditionals.

\bool_lazy_and:nnTF

```

9231 \prg_new_conditional:Npnn \bool_lazy_and:nn #1#2 { p , T , F , TF }
9232 {
9233   \if_predicate:w
9234   \bool_if:nTF {#1} { \bool_if_p:n {#2} } { \c_false_bool }
9235   \prg_return_true:
9236   \else:
9237   \prg_return_false:
9238   \fi:
9239 }

```

(End definition for **\bool_lazy_and:nnTF**. This function is documented on page 110.)

\bool_lazy_any_p:n Go through the list of expressions, stopping whenever an expression is **true**. If the end is reached without finding any **true** expression, then the result is **false**.

\bool_lazy_any:nTF

```

9240 \cs_new:Npn \bool_lazy_any_p:n #1
9241 { \__bool_lazy_any:n #1 \q__bool_recursion_tail \q__bool_recursion_stop }
9242 \prg_new_conditional:Npnn \bool_lazy_any:n #1 { T , F , TF }
9243 {
9244   \if_predicate:w \bool_lazy_any_p:n {#1}
9245   \prg_return_true:
9246   \else:
9247   \prg_return_false:
9248   \fi:
9249 }
9250 \cs_new:Npn \__bool_lazy_any:n #1
9251 {
9252   \__bool_if_recursion_tail_stop_do:nn {#1} { \c_false_bool }
9253   \bool_if:nT {#1}

```

```

9254     { \_bool_use_i_delimit_by_q_recursion_stop:nw { \c_true_bool } }
9255     \_bool_lazy_any:n
9256   }

```

(End definition for \bool_lazy_any:nTF and _bool_lazy_any:n. This function is documented on page 111.)

\bool_lazy_or_p:nn Only evaluate the second expression if the first is false.

```

\bool_lazy_or:nnTF
9257 \prg_new_conditional:Npnn \bool_lazy_or:nn #1#2 { p , T , F , TF }
9258 {
9259   \if_predicate:w
9260     \bool_if:nTF {#1} { \c_true_bool } { \bool_if_p:n {#2} }
9261     \prg_return_true:
9262   \else:
9263     \prg_return_false:
9264   \fi:
9265 }

```

(End definition for \bool_lazy_or:nnTF. This function is documented on page 111.)

\bool_not_p:n The Not variant just reverses the outcome of \bool_if_p:n. Can be optimized but this is nice and simple and according to the implementation plan. Not even particularly useful to have it when the infix notation is easier to use.

```

9266 \cs_new:Npn \bool_not_p:n #1 { \bool_if_p:n { ! ( #1 ) } }

```

(End definition for \bool_not_p:n. This function is documented on page 111.)

\bool_xor_p:nn Exclusive or. If the boolean expressions have same truth value, return false, otherwise return true.

```

\bool_xor:nnTF
9267 \prg_new_conditional:Npnn \bool_xor:nn #1#2 { p , T , F , TF }
9268 {
9269   \bool_if:nT {#1} \reverse_if:N
9270   \if_predicate:w \bool_if_p:n {#2}
9271     \prg_return_true:
9272   \else:
9273     \prg_return_false:
9274   \fi:
9275 }

```

(End definition for \bool_xor:nnTF. This function is documented on page 111.)

14.6 Logical loops

\bool_while_do:Nn A while loop where the boolean is tested before executing the statement. The “while” version executes the code as long as the boolean is true; the “until” version executes the code as long as the boolean is false.

\bool_while_do:cn
\bool_until_do:Nn
\bool_until_do:cn

```

9276 \cs_new:Npn \bool_while_do:Nn #1#2
9277   { \bool_if:NT #1 { #2 \bool_while_do:Nn #1 {#2} } }
9278 \cs_new:Npn \bool_until_do:Nn #1#2
9279   { \bool_if:NF #1 { #2 \bool_until_do:Nn #1 {#2} } }
9280 \cs_generate_variant:Nn \bool_while_do:Nn { c }
9281 \cs_generate_variant:Nn \bool_until_do:Nn { c }

```

(End definition for \bool_while_do:Nn and \bool_until_do:Nn. These functions are documented on page 111.)

`\bool_do_while:Nn` A do-while loop where the body is performed at least once and the boolean is tested after executing the body. Otherwise identical to the above functions.

```

\bool_do_while:cn
\bool_do_until:Nn
\bool_do_until:cn
9282 \cs_new:Npn \bool_do_while:Nn #1#2
9283 { #2 \bool_if:NT #1 { \bool_do_while:Nn #1 {#2} } }
9284 \cs_new:Npn \bool_do_until:Nn #1#2
9285 { #2 \bool_if:NF #1 { \bool_do_until:Nn #1 {#2} } }
9286 \cs_generate_variant:Nn \bool_do_while:Nn { c }
9287 \cs_generate_variant:Nn \bool_do_until:Nn { c }

```

(End definition for `\bool_do_while:Nn` and `\bool_do_until:Nn`. These functions are documented on page 111.)

`\bool_while_do:nn` Loop functions with the test either before or after the first body expansion.

```

\bool_do_while:nn
\bool_until_do:nn
\bool_do_while:nn
\bool_do_until:nn
9288 \cs_new:Npn \bool_while_do:nn #1#2
9289 {
9290   \bool_if:nT {#1}
9291   {
9292     #2
9293     \bool_while_do:nn {#1} {#2}
9294   }
9295 }
9296 \cs_new:Npn \bool_do_while:nn #1#2
9297 {
9298   #2
9299   \bool_if:nT {#1} { \bool_do_while:nn {#1} {#2} }
9300 }
9301 \cs_new:Npn \bool_until_do:nn #1#2
9302 {
9303   \bool_if:nF {#1}
9304   {
9305     #2
9306     \bool_until_do:nn {#1} {#2}
9307   }
9308 }
9309 \cs_new:Npn \bool_do_until:nn #1#2
9310 {
9311   #2
9312   \bool_if:nF {#1} { \bool_do_until:nn {#1} {#2} }
9313 }

```

(End definition for `\bool_while_do:nn` and others. These functions are documented on page 112.)

14.7 Producing multiple copies

```
9314 <@@=prg>
```

`\prg_replicate:nn` This function uses a cascading csname technique by David Kastrup (who else :-)

The idea is to make the input 25 result in first adding five, and then 20 copies of the code to be replicated. The technique uses cascading csnames which means that we start building several csnames so we end up with a list of functions to be called in reverse order. This is important here (and other places) because it means that we can for instance make the function that inserts five copies of something to also hand down ten to the next function in line. This is exactly what happens here: in the example with 25 then the next function is the one that inserts two copies but it sees the ten copies handed

```

\__prg_replicate_0:n
\__prg_replicate_1:n
\__prg_replicate_2:n
\__prg_replicate_3:n
\__prg_replicate_4:n
\__prg_replicate_5:n
\__prg_replicate_6:n
\__prg_replicate_7:n
\__prg_replicate_8:n
\__prg_replicate_9:n
\__prg_replicate_first_0:n
\__prg_replicate_first_1:n
\__prg_replicate_first_2:n

```

down by the previous function. In order to avoid the last function to insert say, 100 copies of the original argument just to gobble them again we define separate functions to be inserted first. These functions also close the expansion of `\exp:w`, which ensures that `\prg_replicate:nn` only requires two steps of expansion.

This function has one flaw though: Since it constantly passes down ten copies of its previous argument it severely affects the main memory once you start demanding hundreds of thousands of copies. Now I don't think this is a real limitation for any ordinary use, and if necessary, it is possible to write `\prg_replicate:nn {1000} { \prg_replicate:nn {1000} {<code>} }`. An alternative approach is to create a string of `m's` with `\exp:w` which can be done with just four macros but that method has its own problems since it can exhaust the string pool. Also, it is considerably slower than what we use here so the few extra csnames are well spent I would say.

```

9315 \cs_new:Npn \prg_replicate:nn #1
9316 {
9317   \exp:w
9318   \exp_after:wN \__prg_replicate_first:N
9319   \int_value:w \int_eval:n {#1}
9320   \cs_end:
9321 }
9322 \cs_new:Npn \__prg_replicate:N #1
9323 { \cs:w \__prg_replicate_#1 :n \__prg_replicate:N }
9324 \cs_new:Npn \__prg_replicate_first:N #1
9325 { \cs:w \__prg_replicate_first_#1 :n \__prg_replicate:N }

```

Then comes all the functions that do the hard work of inserting all the copies. The first function takes `:n` as a parameter.

```

9326 \cs_new:Npn \__prg_replicate_ :n #1 { \cs_end: }
9327 \cs_new:cpn { __prg_replicate_0:n } #1
9328 { \cs_end: {#1#1#1#1#1#1#1#1#1#1} }
9329 \cs_new:cpn { __prg_replicate_1:n } #1
9330 { \cs_end: {#1#1#1#1#1#1#1#1#1#1} #1 }
9331 \cs_new:cpn { __prg_replicate_2:n } #1
9332 { \cs_end: {#1#1#1#1#1#1#1#1#1#1} #1#1 }
9333 \cs_new:cpn { __prg_replicate_3:n } #1
9334 { \cs_end: {#1#1#1#1#1#1#1#1#1#1} #1#1#1 }
9335 \cs_new:cpn { __prg_replicate_4:n } #1
9336 { \cs_end: {#1#1#1#1#1#1#1#1#1#1} #1#1#1#1 }
9337 \cs_new:cpn { __prg_replicate_5:n } #1
9338 { \cs_end: {#1#1#1#1#1#1#1#1#1#1} #1#1#1#1#1 }
9339 \cs_new:cpn { __prg_replicate_6:n } #1
9340 { \cs_end: {#1#1#1#1#1#1#1#1#1#1} #1#1#1#1#1#1 }
9341 \cs_new:cpn { __prg_replicate_7:n } #1
9342 { \cs_end: {#1#1#1#1#1#1#1#1#1#1} #1#1#1#1#1#1#1 }
9343 \cs_new:cpn { __prg_replicate_8:n } #1
9344 { \cs_end: {#1#1#1#1#1#1#1#1#1#1} #1#1#1#1#1#1#1#1 }
9345 \cs_new:cpn { __prg_replicate_9:n } #1
9346 { \cs_end: {#1#1#1#1#1#1#1#1#1#1} #1#1#1#1#1#1#1#1#1 }

```

Users shouldn't ask for something to be replicated once or even not at all but...

```

9347 \cs_new:cpn { __prg_replicate_first_-:n } #1
9348 {
9349   \exp_end:
9350   \__kernel_msg_expandable_error:nn { kernel } { negative-replication }

```

```

9351 }
9352 \cs_new:cpn { __prg_replicate_first_0:n } #1 { \exp_end: }
9353 \cs_new:cpn { __prg_replicate_first_1:n } #1 { \exp_end: #1 }
9354 \cs_new:cpn { __prg_replicate_first_2:n } #1 { \exp_end: #1#1 }
9355 \cs_new:cpn { __prg_replicate_first_3:n } #1 { \exp_end: #1#1#1 }
9356 \cs_new:cpn { __prg_replicate_first_4:n } #1 { \exp_end: #1#1#1#1 }
9357 \cs_new:cpn { __prg_replicate_first_5:n } #1 { \exp_end: #1#1#1#1#1 }
9358 \cs_new:cpn { __prg_replicate_first_6:n } #1 { \exp_end: #1#1#1#1#1#1 }
9359 \cs_new:cpn { __prg_replicate_first_7:n } #1 { \exp_end: #1#1#1#1#1#1#1 }
9360 \cs_new:cpn { __prg_replicate_first_8:n } #1 { \exp_end: #1#1#1#1#1#1#1#1 }
9361 \cs_new:cpn { __prg_replicate_first_9:n } #1
9362 { \exp_end: #1#1#1#1#1#1#1#1#1 }

```

(End definition for `\prg_replicate:nn` and others. This function is documented on page 112.)

14.8 Detecting TeX's mode

`\mode_if_vertical_p:` For testing vertical mode. Strikes me here on the bus with David, that as long as we are just talking about returning true and false states, we can just use the primitive conditionals for this and gobbling the `\exp_end:` in the input stream. However this requires knowledge of the implementation so we keep things nice and clean and use the return statements.

```

9363 \prg_new_conditional:Npnn \mode_if_vertical: { p , T , F , TF }
9364 { \if_mode_vertical: \prg_return_true: \else: \prg_return_false: \fi: }

```

(End definition for `\mode_if_vertical:TF`. This function is documented on page 113.)

`\mode_if_horizontal_p:` For testing horizontal mode.

```

\mode_if_horizontal:TF
9365 \prg_new_conditional:Npnn \mode_if_horizontal: { p , T , F , TF }
9366 { \if_mode_horizontal: \prg_return_true: \else: \prg_return_false: \fi: }

```

(End definition for `\mode_if_horizontal:TF`. This function is documented on page 112.)

`\mode_if_inner_p:` For testing inner mode.

```

\mode_if_inner:TF
9367 \prg_new_conditional:Npnn \mode_if_inner: { p , T , F , TF }
9368 { \if_mode_inner: \prg_return_true: \else: \prg_return_false: \fi: }

```

(End definition for `\mode_if_inner:TF`. This function is documented on page 112.)

`\mode_if_math_p:` For testing math mode. At the beginning of an alignment cell, this should be used only inside a non-expandable function.

```

\mode_if_math:TF
9369 \prg_new_conditional:Npnn \mode_if_math: { p , T , F , TF }
9370 { \if_mode_math: \prg_return_true: \else: \prg_return_false: \fi: }

```

(End definition for `\mode_if_math:TF`. This function is documented on page 112.)

14.9 Internal programming functions

`\group_align_safe_begin:` `\group_align_safe_end:` T_EX’s alignment structures present many problems. As Knuth says himself in *T_EX: The Program*: “It’s sort of a miracle whenever `\halign` or `\valign` work, [...]” One problem relates to commands that internally issues a `\cr` but also peek ahead for the next character for use in, say, an optional argument. If the next token happens to be a `&` with category code 4 we get some sort of weird error message because the underlying `\futurelet` stores the token at the end of the alignment template. This could be a `&_4` giving a message like `! Misplaced \cr.` or even worse: it could be the `\endtemplate` token causing even more trouble! To solve this we have to open a special group so that T_EX still thinks it’s on safe ground but at the same time we don’t want to introduce any brace group that may find its way to the output. The following functions help with this by using code documented only in Appendix D of *The T_EXbook*... We place the `\if_false: { \fi:` part at that place so that the successive expansions of `\group_align_safe_begin/end:` are always brace balanced.

```

9371 \cs_new:Npn \group_align_safe_begin:
9372   { \if_int_compare:w \if_false: { \fi: ‘} = \c_zero_int \fi: }
9373 \cs_new:Npn \group_align_safe_end:
9374   { \if_int_compare:w ‘{ = \c_zero_int } \fi: }

```

(End definition for `\group_align_safe_begin:` and `\group_align_safe_end:`. These functions are documented on page 114.)

`\g__kernel_prg_map_int` A nesting counter for mapping.

```

9375 \int_new:N \g__kernel_prg_map_int

```

(End definition for `\g__kernel_prg_map_int:`.)

`\prg_break_point:Nn` `\prg_map_break:Nn` These are defined in `l3basics`, as they are needed “early”. This is just a reminder that is the case!

(End definition for `\prg_break_point:Nn` and `\prg_map_break:Nn`. These functions are documented on page 113.)

`\prg_break_point:` Also done in `l3basics`.

`\prg_break:`

`\prg_break:n`

(End definition for `\prg_break_point:`, `\prg_break:`, and `\prg_break:n`. These functions are documented on page 114.)

```

9376 \</package>

```

15 l3sys implementation

```

9377 <@@=sys>

```

15.1 Kernel code

```

9378 <*package>

```

```

9379 <*tex>

```

15.1.1 Detecting the engine

`__sys_const:n` Set the T, F, TF, p forms of #1 to be constants equal to the result of evaluating the boolean expression #2.

```

9380 \cs_new_protected:Npn \__sys_const:n #1#2
9381   {

```

```

9382 \bool_if:nTF {#2}
9383 {
9384   \cs_new_eq:cN { #1 :T } \use:n
9385   \cs_new_eq:cN { #1 :F } \use_none:n
9386   \cs_new_eq:cN { #1 :TF } \use_i:nn
9387   \cs_new_eq:cN { #1 _p: } \c_true_bool
9388 }
9389 {
9390   \cs_new_eq:cN { #1 :T } \use_none:n
9391   \cs_new_eq:cN { #1 :F } \use:n
9392   \cs_new_eq:cN { #1 :TF } \use_ii:nn
9393   \cs_new_eq:cN { #1 _p: } \c_false_bool
9394 }
9395 }

```

(End definition for `_sys_const:nn`.)

`\sys_if_engine luatex_p:` Set up the engine tests on the basis exactly one test should be true. Mainly a case of looking for the appropriate marker primitive.

```

\sys_if_engine luatex:TF
\sys_if_engine pdftex_p:
\sys_if_engine pdftex:TF
\sys_if_engine ptex_p:
\sys_if_engine ptex:TF
\sys_if_engine uptex_p:
\sys_if_engine uptex:TF
\sys_if_engine xetex_p:
\sys_if_engine xetex:TF
\c_sys_engine_str
9396 \str_const:Nx \c_sys_engine_str
9397 {
9398   \cs_if_exist:NT \tex luatexversion:D { luatex }
9399   \cs_if_exist:NT \tex pdftexversion:D { pdftex }
9400   \cs_if_exist:NT \tex kanjiskip:D
9401   {
9402     \cs_if_exist:NTF \tex enableecjktoken:D
9403     { uptex }
9404     { ptex }
9405   }
9406   \cs_if_exist:NT \tex XeTeXversion:D { xetex }
9407 }
9408 \tl_map_inline:nn { { luatex } { pdftex } { ptex } { uptex } { xetex } }
9409 {
9410   \_sys_const:nn { sys_if_engine_ #1 }
9411   { \str_if_eq_p:Vn \c_sys_engine_str {#1} }
9412 }

```

(End definition for `\sys_if_engine luatex:TF` and others. These functions are documented on page 115.)

`\c_sys_engine_exec_str` Take the functions defined above, and set up the engine and format names. `\c_sys_engine_exec_str` differs from `\c_sys_engine_str` as it is the *actual* engine name, not a “filtered” version. It differs for `ptex` and `uptex`, which have a leading `e`, and for `luatex`, because `LaTeX` uses the `LuaHBTeX` engine.

`\c_sys_engine_format_str` is quite similar to `\c_sys_engine_str`, except that it differentiates `pdflatex` from `latex` (which is `pdfTeX` in DVI mode). This differentiation, however, is reliable only if the user doesn’t change `\tex_pdfoutput:D` before loading this code.

```

9413 \group_begin:
9414   \cs_set_eq:NN \lua_now:e \tex_directlua:D
9415   \str_const:Nx \c_sys_engine_exec_str
9416   {
9417     \sys_if_engine pdftex:T { pdf }
9418     \sys_if_engine xetex:T { xe }

```



```

9419 \sys_if_engine_ptex:T { ep }
9420 \sys_if_engine_uptex:T { eup }
9421 \sys_if_engine luatex:T
9422 {
9423   lua \lua_now:e
9424   {
9425     if (pcall(require, 'luaharfbuzz')) then ~
9426       tex.print("hb") ~
9427     end
9428   }
9429 }
9430 tex
9431 }
9432 \group_end:
9433 \str_const:Nx \c_sys_engine_format_str
9434 {
9435   \cs_if_exist:NTF \fmtname
9436   {
9437     \bool_lazy_or:nnTF
9438     { \str_if_eq_p:Vn \fmtname { plain } }
9439     { \str_if_eq_p:Vn \fmtname { LaTeX2e } }
9440     {
9441       \sys_if_engine_pdftex:T
9442       { \int_compare:nNnT { \tex_pdfoutput:D } = { 1 } { pdf } }
9443       \sys_if_engine_xetex:T { xe }
9444       \sys_if_engine_ptex:T { p }
9445       \sys_if_engine_uptex:T { up }
9446       \sys_if_engine luatex:T
9447       {
9448         \int_compare:nNnT { \tex_pdfoutput:D } = { 0 } { dvi }
9449         lua
9450       }
9451       \str_if_eq:VnTF \fmtname { LaTeX2e }
9452       { latex }
9453       {
9454         \bool_lazy_and:nnT
9455         { \sys_if_engine_pdftex_p: }
9456         { \int_compare_p:nNn { \tex_pdfoutput:D } = { 0 } }
9457         { e }
9458       }
9459     }
9460   }
9461   { \fmtname }
9462 }
9463 { unknown }
9464 }

```

(End definition for `\c_sys_engine_exec_str` and `\c_sys_engine_format_str`. These variables are documented on page 116.)

15.1.2 Randomness

This candidate function is placed there because `\sys_if_rand_exist:TF` is used in `l3fp-rand`.

`\sys_if_rand_exist_p:` Currently, randomness exists under pdfTeX, LuaTeX, pTeX and upTeX.

`\sys_if_rand_exist:TF`

```

9465 \__sys_const:nn { sys_if_rand_exist }
9466 { \cs_if_exist_p:N \tex_uniformdeviate:D }

```

(End definition for `\sys_if_rand_exist:TF`. This function is documented on page 274.)

15.1.3 Platform

`\sys_if_platform_unix_p:` Setting these up requires the file module (file lookup), so is actually implemented there.

`\sys_if_platform_unix:TF`

`\sys_if_platform_windows_p:` (End definition for `\sys_if_platform_unix:TF`, `\sys_if_platform_windows:TF`, and `\c_sys_platform_str`. These functions are documented on page 116.)

`\sys_if_platform_windows:TF`

`\c_sys_platform_str`

15.1.4 Configurations

`\sys_load_backend:n` Loading the backend code is pretty simply: check that the backend is valid, then load it up.

`__sys_load_backend_check:N`

`\c_sys_backend_str`

```

9467 \cs_new_protected:Npn \sys_load_backend:n #1
9468 {
9469   \sys_finalise:
9470   \str_if_exist:NTF \c_sys_backend_str
9471   {
9472     \str_if_eq:VnF \c_sys_backend_str {#1}
9473     { \__kernel_msg_error:nn { sys } { backend-set } }
9474   }
9475   {
9476     \tl_if_blank:nF {#1}
9477     { \tl_set:Nn \g__sys_backend_tl {#1} }
9478     \__sys_load_backend_check:N \g__sys_backend_tl
9479     \str_const:Nx \c_sys_backend_str { \g__sys_backend_tl }
9480     \__kernel_sys_configuration_load:n
9481     { l3backend- \c_sys_backend_str }
9482   }
9483 }
9484 \cs_new_protected:Npn \__sys_load_backend_check:N #1
9485 {
9486   \sys_if_engine_xetex:TF
9487   {
9488     \str_case:VnF #1
9489     {
9490       { dvisvgm } { }
9491       { xdvipdfmx } { \tl_gset:Nn #1 { xetex } }
9492       { xetex } { }
9493     }
9494     {
9495       \__kernel_msg_error:nxxx { sys } { wrong-backend }
9496       #1 { xetex }
9497       \tl_gset:Nn #1 { xetex }
9498     }
9499   }
9500   {
9501     \sys_if_output_pdf:TF
9502     {
9503       \str_if_eq:VnTF #1 { pdfmode }

```

```

9504         {
9505             \sys_if_engine luatex:TF
9506             { \tl_gset:Nn #1 { luatex } }
9507             { \tl_gset:Nn #1 { pdftex } }
9508         }
9509         {
9510             \bool_lazy_or:nnF
9511             { \str_if_eq_p:Vn #1 { luatex } }
9512             { \str_if_eq_p:Vn #1 { pdftex } }
9513             {
9514                 \__kernel_msg_error:nxxx { sys } { wrong-backend }
9515                 #1 { \sys_if_engine luatex:TF { luatex } { pdftex } }
9516                 \sys_if_engine luatex:TF
9517                 { \tl_gset:Nn #1 { luatex } }
9518                 { \tl_gset:Nn #1 { pdftex } }
9519             }
9520         }
9521     }
9522     {
9523         \str_case:VnF #1
9524         {
9525             { dvipdfmx } { }
9526             { dvips } { }
9527             { dvisvgm } { }
9528         }
9529         {
9530             \__kernel_msg_error:nxxx { sys } { wrong-backend }
9531             #1 { dvips }
9532             \tl_gset:Nn #1 { dvips }
9533         }
9534     }
9535 }
9536 }

```

(End definition for `\sys_load_backend:n`, `__sys_load_backend_check:N`, and `\c_sys_backend_str`. These functions are documented on page 118.)

```

\g__sys_debug_bool
\g__sys_deprecation_bool
9537 \bool_new:N \g__sys_debug_bool
9538 \bool_new:N \g__sys_deprecation_bool

```

(End definition for `\g__sys_debug_bool` and `\g__sys_deprecation_bool`.)

`\sys_load_debug:` Simple.
`\sys_load_deprecation:`

```

9539 \cs_new_protected:Npn \sys_load_debug:
9540 {
9541     \bool_if:NF \g__sys_debug_bool
9542     { \__kernel_sys_configuration_load:n { l3debug } }
9543     \bool_gset_true:N \g__sys_debug_bool
9544 }
9545 \cs_new_protected:Npn \sys_load_deprecation:
9546 {
9547     \bool_if:NF \g__sys_deprecation_bool
9548     { \__kernel_sys_configuration_load:n { l3deprecation } }
9549     \bool_gset_true:N \g__sys_deprecation_bool

```

9550 }

(End definition for \sys_load_debug: and \sys_load_deprecation:. These functions are documented on page 118.)

15.1.5 Access to the shell

\l__sys_internal_tl

9551 \tl_new:N \l__sys_internal_tl

(End definition for \l__sys_internal_tl.)

\c__sys_marker_tl

The same idea as the marker for rescanning token lists.

9552 \tl_const:Nx \c__sys_marker_tl { : \token_to_str:N : }

(End definition for \c__sys_marker_tl.)

\sys_get_shell:nnN

Setting using a shell is at this level just a slightly specialised file operation, with an additional check for quotes, as these are not supported.

\sys_get_shell:nnN

__sys_get:nnN

__sys_get_do:Nw

9553 \cs_new_protected:Npn \sys_get_shell:nnN #1#2#3

9554 {

9555 \sys_get_shell:nnNF {#1} {#2} #3

9556 { \tl_set:Nn #3 { \q_no_value } }

9557 }

9558 \prg_new_protected_conditional:Npnn \sys_get_shell:nnN #1#2#3 { T , F , TF }

9559 {

9560 \sys_if_shell:TF

9561 { \exp_args:No __sys_get:nnN { \tl_to_str:n {#1} } {#2} #3 }

9562 { \prg_return_false: }

9563 }

9564 \cs_new_protected:Npn __sys_get:nnN #1#2#3

9565 {

9566 \tl_if_in:nnTF {#1} { " }

9567 {

9568 __kernel_msg_error:nnx

9569 { kernel } { quote-in-shell } {#1}

9570 \prg_return_false:

9571 }

9572 {

9573 \group_begin:

9574 \if_false: { \fi:

9575 \int_set_eq:NN \tex_tracingnesting:D \c_zero_int

9576 \exp_args:No \tex_everyeof:D { \c__sys_marker_tl }

9577 #2 \scan_stop:

9578 \exp_after:wN __sys_get_do:Nw

9579 \exp_after:wN #3

9580 \exp_after:wN \prg_do_nothing:

9581 \tex_input:D | "#1" \scan_stop:

9582 \if_false: } \fi:

9583 \prg_return_true:

9584 }

9585 }

9586 \exp_args:Nno \use:nn

9587 { \cs_new_protected:Npn __sys_get_do:Nw #1#2 }

9588 { \c__sys_marker_tl }

```

9589 {
9590   \group_end:
9591   \tl_set:No #1 {#2}
9592 }

```

(End definition for `\sys_get_shell:nNTF` and others. These functions are documented on page 117.)

`\c__sys_shell_stream_int` This is not needed for LuaTeX: shell escape there isn't done using a TeX interface.

```

9593 \sys_if_engine luatex:F
9594 { \int_const:Nn \c__sys_shell_stream_int { 18 } }

```

(End definition for `\c__sys_shell_stream_int`.)

`\sys_shell_now:n` Execute commands through shell escape immediately.
`__sys_shell_now:e` For LuaTeX, we use a pseudo-primitive to do the actual work.

```

9595 </tex>
9596 <*lua>
9597 do
9598   local os_exec = os.execute
9599
9600   local function shellescape(cmd)
9601     local status,msg = os_exec(cmd)
9602     if status == nil then
9603       write_nl("log","runsystem(" .. cmd .. ")...(" .. msg .. ")\n")
9604     elseif status == 0 then
9605       write_nl("log","runsystem(" .. cmd .. ")...executed\n")
9606     else
9607       write_nl("log","runsystem(" .. cmd .. ")...failed " .. (msg or "") .. "\n")
9608     end
9609   end
9610   luacmd("__sys_shell_now:e", function()
9611     shellescape(scan_string())
9612   end, "global", "protected")
9613 </lua>
9614 <*tex>
9615 \sys_if_engine luatex:TF
9616 {
9617   \cs_new_protected:Npn \sys_shell_now:n #1
9618   { \__sys_shell_now:e { \exp_not:n {#1} } }
9619 }
9620 {
9621   \cs_new_protected:Npn \sys_shell_now:n #1
9622   { \iow_now:Nn \c__sys_shell_stream_int {#1} }
9623 }
9624 \cs_generate_variant:Nn \sys_shell_now:n { x }
9625 </tex>

```

(End definition for `\sys_shell_now:n` and `__sys_shell_now:e`. This function is documented on page 118.)

`\sys_shell_shipout:n` Execute commands through shell escape at shipout.
`__sys_shell_shipout:e` For LuaTeX, we use the same helper as above but delayed to using a `late_lua` whatsit.

```

9626 <*lua>
9627   local whatsit_id = node.id'whatsit'

```

```

9628 local latelua_sub = node.subtype'late_lua'
9629 local node_new = node.direct.new
9630 local setfield = node.direct.setwhatsitfield or node.direct.setfield
9631 local node_write = node.direct.write
9632
9633 luacmd("__sys_shell_shipout:e", function()
9634     local cmd = scan_string()
9635     local n = node_new(whatsit_id, latelua_sub)
9636     setfield(n, 'data', function() shellescape(cmd) end)
9637     node_write(n)
9638 end, "global", "protected")
9639 end
9640 </lua>
9641 <*tex>
9642 \sys_if_engine luatex:TF
9643 {
9644     \cs_new_protected:Npn \sys_shell_shipout:n #1
9645     { \__sys_shell_shipout:e { \exp_not:n {#1} } }
9646 }
9647 {
9648     \cs_new_protected:Npn \sys_shell_shipout:n #1
9649     { \iow_shipout:Nn \c__sys_shell_stream_int {#1} }
9650 }
9651 \cs_generate_variant:Nn \sys_shell_shipout:n { x }

```

(End definition for `\sys_shell_shipout:n` and `__sys_shell_shipout:e`. This function is documented on page 118.)

15.2 Dynamic (every job) code

```

\sys_everyjob:
\__sys_everyjob:n
\g__sys_everyjob_tl
9652 \cs_new_protected:Npn \sys_everyjob:
9653 {
9654     \tl_use:N \g__sys_everyjob_tl
9655     \tl_gclear:N \g__sys_everyjob_tl
9656 }
9657 \cs_new_protected:Npn \__sys_everyjob:n #1
9658 { \tl_gput_right:Nn \g__sys_everyjob_tl {#1} }
9659 \tl_new:N \g__sys_everyjob_tl

```

(End definition for `\sys_everyjob:`, `__sys_everyjob:n`, and `\g__sys_everyjob_tl`. This function is documented on page ??.)

15.2.1 The name of the job

`\c_sys_jobname_str` Inherited from the L^AT_EX3 name for the primitive. This *has* to be the primitive as it's set in `\everyjob`. If the user does

```
pdflatex \input some-file-name
```

then `\everyjob` is inserted *before* `\jobname` is changed from `texput`, and thus we would have the wrong result.

```

9660 \__sys_everyjob:n
9661 { \cs_new_eq:NN \c_sys_jobname_str \tex_jobname:D }

```

(End definition for `\c_sys_jobname_str`. This variable is documented on page 115.)

15.2.2 Time and date

`\c_sys_minute_int`
`\c_sys_hour_int`
`\c_sys_day_int`
`\c_sys_month_int`
`\c_sys_year_int`

Copies of the information provided by T_EX. There is a lot of defensive code in package mode: someone may have moved the primitives, and they can only be recovered if we have `\primitive` and it is working correctly. For IniT_EX of course that is all redundant but does no harm.

```

9662 \__sys_everyjob:n
9663 {
9664   \group_begin:
9665   \cs_set:Npn \__sys_tmp:w #1
9666   {
9667     \str_if_eq:eeTF { \cs_meaning:N #1 } { \token_to_str:N #1 }
9668     { #1 }
9669     {
9670       \cs_if_exist:NTF \tex_primitive:D
9671       {
9672         \bool_lazy_and:nnTF
9673         { \sys_if_engine_xetex_p: }
9674         {
9675           \int_compare_p:nNn
9676           { \exp_after:wN \use_none:n \tex_XeTeXrevision:D }
9677           < { 99999 }
9678         }
9679         { 0 }
9680         { \tex_primitive:D #1 }
9681       }
9682       { 0 }
9683     }
9684   }
9685   \int_const:Nn \c_sys_minute_int
9686   { \int_mod:nn { \__sys_tmp:w \time } { 60 } }
9687   \int_const:Nn \c_sys_hour_int
9688   { \int_div_truncate:nn { \__sys_tmp:w \time } { 60 } }
9689   \int_const:Nn \c_sys_day_int { \__sys_tmp:w \day }
9690   \int_const:Nn \c_sys_month_int { \__sys_tmp:w \month }
9691   \int_const:Nn \c_sys_year_int { \__sys_tmp:w \year }
9692   \group_end:
9693 }

```

(End definition for `\c_sys_minute_int` and others. These variables are documented on page 115.)

15.2.3 Random numbers

`\sys_rand_seed:` Unpack the primitive. When random numbers are not available, we return zero after an error (and incidentally make sure the number of expansions needed is the same as with random numbers available).

```

9694 \__sys_everyjob:n
9695 {
9696   \sys_if_rand_exist:TF
9697   { \cs_new:Npn \sys_rand_seed: { \tex_the:D \tex_randomseed:D } }
9698   {
9699     \cs_new:Npn \sys_rand_seed:
9700     {
9701       \int_value:w

```

```

9702         \__kernel_msg_expandable_error:nnn { kernel } { fp-no-random }
9703         { \sys_rand_seed: }
9704         \c_zero_int
9705     }
9706 }
9707 }

```

(End definition for `\sys_rand_seed`:. This function is documented on page 116.)

`\sys_gset_rand_seed:n` The primitive always assigns the seed globally.

```

9708 \__sys_everyjob:n
9709 {
9710     \sys_if_rand_exist:TF
9711     {
9712         \cs_new_protected:Npn \sys_gset_rand_seed:n #1
9713         { \tex_setrandomseed:D \int_eval:n {#1} \exp_stop_f: }
9714     }
9715     {
9716         \cs_new_protected:Npn \sys_gset_rand_seed:n #1
9717         {
9718             \__kernel_msg_error:nnn { kernel } { fp-no-random }
9719             { \sys_gset_rand_seed:n {#1} }
9720         }
9721     }
9722 }

```

(End definition for `\sys_gset_rand_seed:n`. This function is documented on page 117.)

15.2.4 Access to the shell

`\c_sys_shell_escape_int` Expose the engine's shell escape status to the user.

```

9723 \__sys_everyjob:n
9724 {
9725     \int_const:Nn \c_sys_shell_escape_int
9726     {
9727         \sys_if_engine luatex:TF
9728         {
9729             \tex_directlua:D
9730             { \tex_sprint(status.shell_escape~or~os.execute()) }
9731         }
9732         { \tex_shellescape:D }
9733     }
9734 }

```

(End definition for `\c_sys_shell_escape_int`. This variable is documented on page 117.)

`\sys_if_shell_p:` Performs a check for whether shell escape is enabled. The first set of functions returns true if either of restricted or unrestricted shell escape is enabled, while the other two sets of functions return true in only one of these two cases.

`\sys_if_shell:TF`

`\sys_if_shell_unrestricted_p:`

`\sys_if_shell_unrestricted:TF`

`\sys_if_shell_restricted_p:`

`\sys_if_shell_restricted:TF`

```

9735 \__sys_everyjob:n
9736 {
9737     \__sys_const:nn { sys_if_shell }
9738     { \int_compare_p:nNn \c_sys_shell_escape_int > 0 }
9739     \__sys_const:nn { sys_if_shell_unrestricted }

```



```

9740     { \int_compare_p:nNn \c_sys_shell_escape_int = 1 }
9741   \__sys_const:nn { sys_if_shell_restricted }
9742     { \int_compare_p:nNn \c_sys_shell_escape_int = 2 }
9743   }

```

(End definition for `\sys_if_shell:TF`, `\sys_if_shell_unrestricted:TF`, and `\sys_if_shell_restricted:TF`. These functions are documented on page 117.)

15.2.5 Held over from l3file

`\g_file_curr_name_str` See comments about `\c_sys_jobname_str`: here, as soon as there is file input/output, things get “tided up”.

```

9744 \__sys_everyjob:n
9745   { \cs_gset_eq:NN \g_file_curr_name_str \tex_jobname:D }

```

(End definition for `\g_file_curr_name_str`. This variable is documented on page 167.)

15.3 Last-minute code

`\sys_finalise:` A simple hook to finalise the system-dependent layer. This is forced by the backend loader, which is forced by the main loader, so we do not need to include that here.

```

\__sys_finalise:n
\g__sys_finalise_tl
9746 \cs_new_protected:Npn \sys_finalise:
9747   {
9748     \sys_everyjob:
9749     \tl_use:N \g__sys_finalise_tl
9750     \tl_gclear:N \g__sys_finalise_tl
9751   }
9752 \cs_new_protected:Npn \__sys_finalise:n #1
9753   { \tl_gput_right:Nn \g__sys_finalise_tl {#1} }
9754 \tl_new:N \g__sys_finalise_tl

```

(End definition for `\sys_finalise:`, `__sys_finalise:n`, and `\g__sys_finalise_tl`. This function is documented on page 118.)

15.3.1 Detecting the output

`\sys_if_output_dvi_p:` This is a simple enough concept: the two views here are complementary.

```

\sys_if_output_dvi:TF
\sys_if_output_pdf_p:
\sys_if_output_pdf:TF
\c_sys_output_str
9755 \__sys_finalise:n
9756   {
9757     \str_const:Nx \c_sys_output_str
9758     {
9759       \int_compare:nNnTF
9760         { \cs_if_exist_use:NF \tex_pdfoutput:D { 0 } } > { 0 }
9761         { pdf }
9762         { dvi }
9763     }
9764     \__sys_const:nn { sys_if_output_dvi }
9765     { \str_if_eq_p:Vn \c_sys_output_str { dvi } }
9766     \__sys_const:nn { sys_if_output_pdf }
9767     { \str_if_eq_p:Vn \c_sys_output_str { pdf } }
9768   }

```

(End definition for `\sys_if_output_dvi:TF`, `\sys_if_output_pdf:TF`, and `\c_sys_output_str`. These functions are documented on page 116.)

15.3.2 Configurations

`\g__sys_backend_tl` As the backend has to be checked and possibly adjusted, the approach here is to create a variable and use that in a one-shot to set a constant.

```

9769 \tl_new:N \g__sys_backend_tl
9770 \__sys_finalise:n
9771 {
9772   \__kernel_tl_gset:Nx \g__sys_backend_tl
9773   {
9774     \sys_if_engine_xetex:TF
9775     { xetex }
9776     {
9777       \sys_if_output_pdf:TF
9778       {
9779         \sys_if_engine_pdftex:TF
9780         { pdftex }
9781         { luatex }
9782       }
9783       { dvips }
9784     }
9785   }
9786 }
```

If there is a class option set, and recognised, we pick it up: these will over-ride anything set automatically but will themselves be over-written if there is a package option.

```

9787 \__sys_finalise:n
9788 {
9789   \cs_if_exist:NT \@classoptionslist
9790   {
9791     \cs_if_eq:NMF \@classoptionslist \scan_stop:
9792     {
9793       \clist_map_inline:Nn \@classoptionslist
9794       {
9795         \str_case:nnT {#1}
9796         {
9797           { dvipdfmx }
9798           { \tl_gset:Nn \g__sys_backend_tl { dvipdfmx } }
9799           { dvips }
9800           { \tl_gset:Nn \g__sys_backend_tl { dvips } }
9801           { dvisvgm }
9802           { \tl_gset:Nn \g__sys_backend_tl { dvisvgm } }
9803           { pdftex }
9804           { \tl_gset:Nn \g__sys_backend_tl { pdfmode } }
9805           { xetex }
9806           { \tl_gset:Nn \g__sys_backend_tl { xdvipdfmx } }
9807         }
9808         { \clist_remove_all:Nn \@unusedoptionlist {#1} }
9809       }
9810     }
9811   }
9812 }
```

(End definition for `\g__sys_backend_tl`.)

```

9813 </tex>
9814 </package>
```

16 l3clist implementation

The following test files are used for this code: *m3clist002*.

```
9815 \*package
```

```
9816 \@@=clist
```

\c_empty_clist An empty comma list is simply an empty token list.

```
9817 \cs_new_eq:NN \c_empty_clist \c_empty_tl
```

(End definition for `\c_empty_clist`. This variable is documented on page 128.)

\l__clist_internal_clist Scratch space for various internal uses. This comma list variable cannot be declared as such because it comes before `\clist_new:N`

```
9818 \tl_new:N \l__clist_internal_clist
```

(End definition for `\l__clist_internal_clist`.)

\s__clist_mark Internal scan marks.

```
\s__clist_stop 9819 \scan_new:N \s__clist_mark
```

```
9820 \scan_new:N \s__clist_stop
```

(End definition for `\s__clist_mark` and `\s__clist_stop`.)

__clist_use_none_delimit_by_s_stop:w Functions to gobble up to a scan mark.

```
\__clist_use_i_delimit_by_s_stop:nw 9821 \cs_new:Npn \__clist_use_none_delimit_by_s_stop:w #1 \s__clist_stop { }
```

```
9822 \cs_new:Npn \__clist_use_i_delimit_by_s_stop:nw #1 #2 \s__clist_stop {#1}
```

(End definition for `__clist_use_none_delimit_by_s_stop:w` and `__clist_use_i_delimit_by_s_stop:nw`.)

\q__clist_recursion_tail Internal recursion quarks.

```
\q__clist_recursion_stop 9823 \quark_new:N \q__clist_recursion_tail
```

```
9824 \quark_new:N \q__clist_recursion_stop
```

(End definition for `\q__clist_recursion_tail` and `\q__clist_recursion_stop`.)

__clist_if_recursion_tail_break:nN Functions to query recursion quarks.

```
\__clist_if_recursion_tail_stop:n 9825 \__kernel_quark_new_test:N \__clist_if_recursion_tail_break:nN
```

```
9826 \__kernel_quark_new_test:N \__clist_if_recursion_tail_stop:n
```

(End definition for `__clist_if_recursion_tail_break:nN` and `__clist_if_recursion_tail_stop:n`.)

__clist_tmp:w A temporary function for various purposes.

```
9827 \cs_new_protected:Npn \__clist_tmp:w { }
```

(End definition for `__clist_tmp:w`.)

16.1 Removing spaces around items

`__clist_trim_next:w` Called as `\exp:w __clist_trim_next:w \prg_do_nothing: <comma list> ...` it expands to `{<trimmed item>}` where the `<trimmed item>` is the first non-empty result from removing spaces from both ends of comma-delimited items in the `<comma list>`. The `\prg_do_nothing:` marker avoids losing braces. The test for blank items is a somewhat optimized `\tl_if_empty:oTF` construction; if blank, another item is sought, otherwise trim spaces.

```

9828 \cs_new:Npn \__clist_trim_next:w #1 ,
9829 {
9830     \tl_if_empty:oTF { \use_none:nn #1 ? }
9831     { \__clist_trim_next:w \prg_do_nothing: }
9832     { \tl_trim_spaces_apply:oN {#1} \exp_end: }
9833 }
```

(End definition for `__clist_trim_next:w`.)

`__clist_sanitize:n` The auxiliary `__clist_sanitize:Nn` receives a delimiter (`\c_empty_tl` the first time, afterwards a comma) and that item as arguments. Unless we are done with the loop it calls `__clist_wrap_item:w` to unbrace the item (using a comma delimiter is safe since `#2` came from removing spaces from an argument delimited by a comma) and possibly re-brace it if needed.

```

9834 \cs_new:Npn \__clist_sanitize:n #1
9835 {
9836     \exp_after:wN \__clist_sanitize:Nn \exp_after:wN \c_empty_tl
9837     \exp:w \__clist_trim_next:w \prg_do_nothing:
9838     #1 , \q__clist_recursion_tail , \q__clist_recursion_stop
9839 }
9840 \cs_new:Npn \__clist_sanitize:Nn #1#2
9841 {
9842     \__clist_if_recursion_tail_stop:n {#2}
9843     #1 \__clist_wrap_item:w #2 ,
9844     \exp_after:wN \__clist_sanitize:Nn \exp_after:wN ,
9845     \exp:w \__clist_trim_next:w \prg_do_nothing:
9846 }
```

(End definition for `__clist_sanitize:n` and `__clist_sanitize:Nn`.)

`__clist_if_wrap:nTF` True if the argument must be wrapped to avoid getting altered by some clist operations.
`__clist_if_wrap:w` That is the case whenever the argument

- starts or end with a space or contains a comma,
- is empty, or
- consists of a single braced group.

All `l3clist` functions go through the same test when they need to determine whether to brace an item, so it is not a problem that this test has false positives such as “`\s__clist_mark ?`”. If the argument starts or end with a space or contains a comma then one of the three arguments of `__clist_if_wrap:w` will have its end delimiter (partly) in one of the three copies of `#1` in `__clist_if_wrap:nTF`; this has a knock-on effect meaning that the result of the expansion is not empty; in that case, wrap. Otherwise,

the argument is safe unless it starts with a brace group (or is empty) and it is empty or consists of a single n-type argument.

```

9847 \prg_new_conditional:Npnn \__clist_if_wrap:n #1 { TF }
9848 {
9849   \tl_if_empty:oTF
9850   {
9851     \__clist_if_wrap:w
9852     \s__clist_mark ? #1 ~ \s__clist_mark ? ~ #1
9853     \s__clist_mark , ~ \s__clist_mark #1 ,
9854   }
9855   {
9856     \tl_if_head_is_group:nTF { #1 { } }
9857     {
9858       \tl_if_empty:nTF {#1}
9859       { \prg_return_true: }
9860       {
9861         \tl_if_empty:oTF { \use_none:n #1}
9862         { \prg_return_true: }
9863         { \prg_return_false: }
9864       }
9865     }
9866     { \prg_return_false: }
9867   }
9868   { \prg_return_true: }
9869 }
9870 \cs_new:Npn \__clist_if_wrap:w #1 \s__clist_mark ? ~ #2 ~ \s__clist_mark #3 , { }

```

(End definition for __clist_if_wrap:nTF and __clist_if_wrap:w.)

__clist_wrap_item:w Safe items are put in \exp_not:n, otherwise we put an extra set of braces.

```

9871 \cs_new:Npn \__clist_wrap_item:w #1 ,
9872 { \__clist_if_wrap:nTF {#1} { \exp_not:n { {#1} } } { \exp_not:n {#1} } }

```

(End definition for __clist_wrap_item:w.)

16.2 Allocation and initialisation

\clist_new:N Internally, comma lists are just token lists.

```

\clist_new:c
9873 \cs_new_eq:NN \clist_new:N \tl_new:N
9874 \cs_new_eq:NN \clist_new:c \tl_new:c

```

(End definition for \clist_new:N. This function is documented on page 119.)

\clist_const:Nn Creating and initializing a constant comma list is done by sanitizing all items (stripping spaces and braces).

```

\clist_const:cn
\clist_const:Nx
\clist_const:cx
9875 \cs_new_protected:Npn \clist_const:Nn #1#2
9876 { \tl_const:Nx #1 { \__clist_sanitize:n {#2} } }
9877 \cs_generate_variant:Nn \clist_const:Nn { c , Nx , cx }

```

(End definition for \clist_const:Nn. This function is documented on page 120.)

`\clist_clear:N` Clearing comma lists is just the same as clearing token lists.

```
\clist_clear:c      9878 \cs_new_eq:NN \clist_clear:N \tl_clear:N
\clist_gclear:N     9879 \cs_new_eq:NN \clist_clear:c \tl_clear:c
\clist_gclear:c     9880 \cs_new_eq:NN \clist_gclear:N \tl_gclear:N
                   9881 \cs_new_eq:NN \clist_gclear:c \tl_gclear:c
```

(End definition for `\clist_clear:N` and `\clist_gclear:N`. These functions are documented on page 120.)

`\clist_clear_new:N` Once again a copy from the token list functions.

```
\clist_clear_new:c  9882 \cs_new_eq:NN \clist_clear_new:N \tl_clear_new:N
\clist_gclear_new:N 9883 \cs_new_eq:NN \clist_clear_new:c \tl_clear_new:c
\clist_gclear_new:c 9884 \cs_new_eq:NN \clist_gclear_new:N \tl_gclear_new:N
                   9885 \cs_new_eq:NN \clist_gclear_new:c \tl_gclear_new:c
```

(End definition for `\clist_clear_new:N` and `\clist_gclear_new:N`. These functions are documented on page 120.)

`\clist_set_eq:NN` Once again, these are simple copies from the token list functions.

```
\clist_set_eq:cN    9886 \cs_new_eq:NN \clist_set_eq:NN \tl_set_eq:NN
\clist_set_eq:Nc    9887 \cs_new_eq:NN \clist_set_eq:Nc \tl_set_eq:Nc
\clist_set_eq:cc    9888 \cs_new_eq:NN \clist_set_eq:cN \tl_set_eq:cN
\clist_gset_eq:NN   9889 \cs_new_eq:NN \clist_set_eq:cc \tl_set_eq:cc
\clist_gset_eq:cN   9890 \cs_new_eq:NN \clist_gset_eq:NN \tl_gset_eq:NN
\clist_gset_eq:Nc   9891 \cs_new_eq:NN \clist_gset_eq:Nc \tl_gset_eq:Nc
\clist_gset_eq:cN   9892 \cs_new_eq:NN \clist_gset_eq:cN \tl_gset_eq:cN
\clist_gset_eq:cc   9893 \cs_new_eq:NN \clist_gset_eq:cc \tl_gset_eq:cc
```

(End definition for `\clist_set_eq:NN` and `\clist_gset_eq:NN`. These functions are documented on page 120.)

`\clist_set_from_seq:NN` Setting a comma list from a comma-separated list is done using a simple mapping. Safe items are put in `\exp_not:n`, otherwise we put an extra set of braces. The first comma must be removed, except in the case of an empty comma-list.

```
\clist_set_from_seq:cN 9894 \cs_new_protected:Npn \clist_set_from_seq:NN
\clist_set_from_seq:Nc 9895 { \__clist_set_from_seq:NNNN \clist_clear:N \__kernel_tl_set:Nx }
\clist_set_from_seq:cc 9896 \cs_new_protected:Npn \clist_gset_from_seq:NN
\clist_gset_from_seq:cN 9897 { \__clist_set_from_seq:NNNN \clist_gclear:N \__kernel_tl_gset:Nx }
\clist_gset_from_seq:Nc 9898 \cs_new_protected:Npn \__clist_set_from_seq:NNNN #1#2#3#4
\clist_gset_from_seq:cc 9899 {
\__clist_set_from_seq:NNNN 9900 \seq_if_empty:NTF #4
\__clist_set_from_seq:n    9901 { #1 #3 }
                           9902 {
                           9903   #2 #3
                           9904   {
                           9905     \exp_after:wN \use_none:n \exp:w \exp_end_continue_f:w
                           9906     \seq_map_function:NN #4 \__clist_set_from_seq:n
                           9907   }
                           9908 }
                           9909 }
                           9910 \cs_new:Npn \__clist_set_from_seq:n #1
                           9911 {
                           9912   ,
                           9913   \__clist_if_wrap:NTF {#1}
                           9914   { \exp_not:n { {#1} } } }
```

```

9915     { \exp_not:n {#1} }
9916   }
9917   \cs_generate_variant:Nn \clist_set_from_seq:NN { Nc }
9918   \cs_generate_variant:Nn \clist_set_from_seq:NN { c , cc }
9919   \cs_generate_variant:Nn \clist_gset_from_seq:NN { Nc }
9920   \cs_generate_variant:Nn \clist_gset_from_seq:NN { c , cc }

```

(End definition for \clist_set_from_seq:NN and others. These functions are documented on page 120.)

```

\clist_concat:NNN Concatenating comma lists is not quite as easy as it seems, as there needs to be the
\clist_concat:ccc correct addition of a comma to the output. So a little work to do.
\clist_gconcat:NNN
\clist_gconcat:ccc
\__clist_concat:NNNN
9921 \cs_new_protected:Npn \clist_concat:NNN
9922 { \__clist_concat:NNNN \__kernel_tl_set:Nx }
9923 \cs_new_protected:Npn \clist_gconcat:NNN
9924 { \__clist_concat:NNNN \__kernel_tl_gset:Nx }
9925 \cs_new_protected:Npn \__clist_concat:NNNN #1#2#3#4
9926 {
9927   #1 #2
9928   {
9929     \exp_not:o #3
9930     \clist_if_empty:NF #3 { \clist_if_empty:NF #4 { , } }
9931     \exp_not:o #4
9932   }
9933 }
9934 \cs_generate_variant:Nn \clist_concat:NNN { ccc }
9935 \cs_generate_variant:Nn \clist_gconcat:NNN { ccc }

```

(End definition for \clist_concat:NNN, \clist_gconcat:NNN, and __clist_concat:NNNN. These functions are documented on page 120.)

```

\clist_if_exist_p:N Copies of the cs functions defined in l3basics.
\clist_if_exist_p:c
\clist_if_exist:NTF
\clist_if_exist:cTF
9936 \prg_new_eq_conditional:NNn \clist_if_exist:N \cs_if_exist:N
9937 { TF , T , F , p }
9938 \prg_new_eq_conditional:NNn \clist_if_exist:c \cs_if_exist:c
9939 { TF , T , F , p }

```

(End definition for \clist_if_exist:NTF. This function is documented on page 120.)

16.3 Adding data to comma lists

```

\clist_set:Nn
\clist_set:NV
\clist_set:No
\clist_set:Nx
\clist_set:cn
\clist_set:cV
\clist_set:co
\clist_set:cx
\clist_gset:Nn
\clist_gset:NV
\clist_gset:No
\clist_gset:Nx
\clist_gset:cn
\clist_gset:cV
\clist_gset:co
\clist_gset:cx
\clist_put_left:cx
\clist_gput_left:Nn
\clist_gput_left:NV
\clist_gput_left:No
\clist_gput_left:Nx
\clist_gput_left:cn
\clist_gput_left:cV
\clist_gput_left:co

```

(End definition for \clist_set:Nn and \clist_gset:Nn. These functions are documented on page 121.)

Everything is based on concatenation after storing in \l__clist_internal_clist. This avoids having to worry here about space-trimming and so on.

```

9946 \cs_new_protected:Npn \clist_put_left:Nn
9947 { \__clist_put_left:NNNN \clist_concat:NNN \clist_set:Nn }
9948 \cs_new_protected:Npn \clist_gput_left:Nn

```

```

9949 { \_clist_put_left:NNN \clist_gconcat:NNN \clist_set:Nn }
9950 \cs_new_protected:Npn \_clist_put_left:NNNn #1#2#3#4
9951 {
9952   #2 \l__clist_internal_clist {#4}
9953   #1 #3 \l__clist_internal_clist #3
9954 }
9955 \cs_generate_variant:Nn \clist_put_left:Nn { NV , No , Nx }
9956 \cs_generate_variant:Nn \clist_put_left:Nn { c , cV , co , cx }
9957 \cs_generate_variant:Nn \clist_gput_left:Nn { NV , No , Nx }
9958 \cs_generate_variant:Nn \clist_gput_left:Nn { c , cV , co , cx }

```

(End definition for `\clist_put_left:Nn`, `\clist_gput_left:Nn`, and `_clist_put_left:NNNn`. These functions are documented on page 121.)

```

\clist_put_right:Nn
\clist_put_right:NV
\clist_put_right:No
\clist_put_right:Nx
\clist_put_right:cn
\clist_put_right:cV
\clist_put_right:co
\clist_put_right:cx
\clist_gput_right:Nn
\clist_gput_right:NV
\clist_gput_right:No
\clist_gput_right:Nx
\clist_gput_right:cn
\clist_gput_right:cV
\clist_gput_right:co
\clist_gput_right:cx
\_clist_put_right:NNNn

```

```

9959 \cs_new_protected:Npn \clist_put_right:Nn
9960 { \_clist_put_right:NNNn \clist_concat:NNN \clist_set:Nn }
9961 \cs_new_protected:Npn \clist_gput_right:Nn
9962 { \_clist_put_right:NNNn \clist_gconcat:NNN \clist_set:Nn }
9963 \cs_new_protected:Npn \_clist_put_right:NNNn #1#2#3#4
9964 {
9965   #2 \l__clist_internal_clist {#4}
9966   #1 #3 #3 \l__clist_internal_clist
9967 }
9968 \cs_generate_variant:Nn \clist_put_right:Nn { NV , No , Nx }
9969 \cs_generate_variant:Nn \clist_put_right:Nn { c , cV , co , cx }
9970 \cs_generate_variant:Nn \clist_gput_right:Nn { NV , No , Nx }
9971 \cs_generate_variant:Nn \clist_gput_right:Nn { c , cV , co , cx }

```

(End definition for `\clist_put_right:Nn`, `\clist_gput_right:Nn`, and `_clist_put_right:NNNn`. These functions are documented on page 121.)

16.4 Comma lists as stacks

`\clist_get:NN` Getting an item from the left of a comma list is pretty easy: just trim off the first item using the comma. No need to trim spaces as comma-list *variables* are assumed to have “cleaned-up” items. (Note that grabbing a comma-delimited item removes an outer pair of braces if present, exactly as needed to uncover the underlying item.)

```

9972 \cs_new_protected:Npn \clist_get:NN #1#2
9973 {
9974   \if_meaning:w #1 \c_empty_clist
9975     \tl_set:Nn #2 { \q_no_value }
9976   \else:
9977     \exp_after:wN \_clist_get:wN #1 , \s__clist_stop #2
9978   \fi:
9979 }
9980 \cs_new_protected:Npn \_clist_get:wN #1 , #2 \s__clist_stop #3
9981 { \tl_set:Nn #3 {#1} }
9982 \cs_generate_variant:Nn \clist_get:NN { c }

```

(End definition for `\clist_get:NN` and `_clist_get:wN`. This function is documented on page 126.)

`\clist_pop:NN` An empty clist leads to `\q_no_value`, otherwise grab until the first comma and assign to the variable. The second argument of `_clist_pop:wwNNN` is a comma list ending in a

```

\clist_pop:NN
\clist_pop:cn
\clist_gpop:NN
\clist_gpop:cn
\_clist_pop:NNN
\_clist_pop:wwNNN
\_clist_pop:wN

```


comma and `\s__clist_mark`, unless the original clist contained exactly one item: then the argument is just `\s__clist_mark`. The next auxiliary picks either `\exp_not:n` or `\use_none:n` as #2, ensuring that the result can safely be an empty comma list.

```

9983 \cs_new_protected:Npn \clist_pop:NN
9984 { \__clist_pop:NNN \__kernel_tl_set:Nx }
9985 \cs_new_protected:Npn \clist_gpop:NN
9986 { \__clist_pop:NNN \__kernel_tl_gset:Nx }
9987 \cs_new_protected:Npn \__clist_pop:NNN #1#2#3
9988 {
9989   \if_meaning:w #2 \c_empty_clist
9990     \tl_set:Nn #3 { \q_no_value }
9991   \else:
9992     \exp_after:wN \__clist_pop:wwNNN #2 , \s__clist_mark \s__clist_stop #1#2#3
9993   \fi:
9994 }
9995 \cs_new_protected:Npn \__clist_pop:wwNNN #1 , #2 \s__clist_stop #3#4#5
9996 {
9997   \tl_set:Nn #5 {#1}
9998   #3 #4
9999   {
10000     \__clist_pop:wN \prg_do_nothing:
10001       #2 \exp_not:o
10002       , \s__clist_mark \use_none:n
10003     \s__clist_stop
10004   }
10005 }
10006 \cs_new:Npn \__clist_pop:wN #1 , \s__clist_mark #2 #3 \s__clist_stop { #2 {#1} }
10007 \cs_generate_variant:Nn \clist_pop:NN { c }
10008 \cs_generate_variant:Nn \clist_gpop:NN { c }

```

(End definition for `\clist_pop:NN` and others. These functions are documented on page 126.)

```

\clist_get:NNTF The same, as branching code: very similar to the above.
\clist_get:cNTF 10009 \prg_new_protected_conditional:Npnn \clist_get:NN #1#2 { T , F , TF }
\clist_pop:NNTF 10010 {
\clist_pop:cNTF 10011   \if_meaning:w #1 \c_empty_clist
\clist_gpop:NNTF 10012   \prg_return_false:
\clist_gpop:cNTF 10013   \else:
\__clist_pop_TF:NNN 10014     \exp_after:wN \__clist_get:wN #1 , \s__clist_stop #2
10015     \prg_return_true:
10016   \fi:
10017 }
10018 \prg_generate_conditional_variant:Nnn \clist_get:NN { c } { T , F , TF }
10019 \prg_new_protected_conditional:Npnn \clist_pop:NN #1#2 { T , F , TF }
10020 { \__clist_pop_TF:NNN \__kernel_tl_set:Nx #1 #2 }
10021 \prg_new_protected_conditional:Npnn \clist_gpop:NN #1#2 { T , F , TF }
10022 { \__clist_pop_TF:NNN \__kernel_tl_gset:Nx #1 #2 }
10023 \cs_new_protected:Npn \__clist_pop_TF:NNN #1#2#3
10024 {
10025   \if_meaning:w #2 \c_empty_clist
10026     \prg_return_false:
10027   \else:
10028     \exp_after:wN \__clist_pop:wwNNN #2 , \s__clist_mark \s__clist_stop #1#2#3
10029     \prg_return_true:

```

```

10030     \fi:
10031   }
10032   \prg_generate_conditional_variant:Nnn \clist_pop:NN { c } { T , F , TF }
10033   \prg_generate_conditional_variant:Nnn \clist_gpop:NN { c } { T , F , TF }

```

(End definition for `\clist_get:NNTF` and others. These functions are documented on page 126.)

`\clist_push:Nn` Pushing to a comma list is the same as adding on the left.

```

\clist_push:NV 10034 \cs_new_eq:NN \clist_push:Nn \clist_put_left:Nn
\clist_push:No 10035 \cs_new_eq:NN \clist_push:NV \clist_put_left:NV
\clist_push:Nx 10036 \cs_new_eq:NN \clist_push:No \clist_put_left:No
\clist_push:cn 10037 \cs_new_eq:NN \clist_push:Nx \clist_put_left:Nx
\clist_push:cV 10038 \cs_new_eq:NN \clist_push:cn \clist_put_left:cn
\clist_push:co 10039 \cs_new_eq:NN \clist_push:cV \clist_put_left:cV
\clist_push:cx 10040 \cs_new_eq:NN \clist_push:co \clist_put_left:co
\clist_gpush:Nn 10041 \cs_new_eq:NN \clist_gpush:Nn \clist_gput_left:Nn
\clist_gpush:NV 10042 \cs_new_eq:NN \clist_gpush:NV \clist_gput_left:NV
\clist_gpush:No 10043 \cs_new_eq:NN \clist_gpush:No \clist_gput_left:No
\clist_gpush:Nx 10044 \cs_new_eq:NN \clist_gpush:Nx \clist_gput_left:Nx
\clist_gpush:cn 10045 \cs_new_eq:NN \clist_gpush:cn \clist_gput_left:cn
\clist_gpush:cV 10046 \cs_new_eq:NN \clist_gpush:cV \clist_gput_left:cV
\clist_gpush:co 10047 \cs_new_eq:NN \clist_gpush:co \clist_gput_left:co
\clist_gpush:cx 10048 \cs_new_eq:NN \clist_gpush:cx \clist_gput_left:cx

```

(End definition for `\clist_push:Nn` and `\clist_gpush:Nn`. These functions are documented on page 127.)

16.5 Modifying comma lists

`\l__clist_internal_remove_clist` An internal comma list and a sequence for the removal routines.
`\l__clist_internal_remove_seq`

```

10050 \clist_new:N \l__clist_internal_remove_clist
10051 \seq_new:N \l__clist_internal_remove_seq

```

(End definition for `\l__clist_internal_remove_clist` and `\l__clist_internal_remove_seq`.)

`\clist_remove_duplicates:N` Removing duplicates means making a new list then copying it.

```

\clist_remove_duplicates:c 10052 \cs_new_protected:Npn \clist_remove_duplicates:N
\clist_gremove_duplicates:N 10053 { \__clist_remove_duplicates:NN \clist_set_eq:NN }
\clist_gremove_duplicates:c 10054 \cs_new_protected:Npn \clist_gremove_duplicates:N
  \__clist_remove_duplicates:NN 10055 { \__clist_remove_duplicates:NN \clist_gset_eq:NN }
  \__clist_remove_duplicates:NN 10056 \cs_new_protected:Npn \__clist_remove_duplicates:NN #1#2
  {
    10057 {
      10058 \clist_clear:N \l__clist_internal_remove_clist
      10059 \clist_map_inline:Nn #2
      10060 {
        10061 \clist_if_in:NnF \l__clist_internal_remove_clist {##1}
        10062 { \clist_put_right:Nn \l__clist_internal_remove_clist {##1} }
      }
      10063 }
      10064 #1 #2 \l__clist_internal_remove_clist
      10065 }
    10066 \cs_generate_variant:Nn \clist_remove_duplicates:N { c }
    10067 \cs_generate_variant:Nn \clist_gremove_duplicates:N { c }

```

(End definition for `\clist_remove_duplicates:N`, `\clist_gremove_duplicates:N`, and `__clist_remove_duplicates:NN`. These functions are documented on page 122.)

```

\clist_remove_all:Nn
\clist_remove_all:cn
\clist_gremove_all:Nn
\clist_gremove_all:cn
\__clist_remove_all:NNNn
\__clist_remove_all:w
\__clist_remove_all:

```

The method used here for safe items is very similar to `\tl_replace_all:Nnn`. However, if the item contains commas or leading/trailing spaces, or is empty, or consists of a single brace group, we know that it can only appear within braces so the code would fail; instead just convert to a sequence and do the removal with `l3seq` code (it involves somewhat elaborate code to do most of the work expandably but the final token list comparisons non-expandably).

For “safe” items, build a function delimited by the $\langle item \rangle$ that should be removed, surrounded with commas, and call that function followed by the expanded comma list, and another copy of the $\langle item \rangle$. The loop is controlled by the argument grabbed by `__clist_remove_all:w`: when the item was found, the `\s__clist_mark` delimiter used is the one inserted by `__clist_tmp:w`, and `__clist_use_none_delimit_by_s_stop:w` is deleted. At the end, the final $\langle item \rangle$ is grabbed, and the argument of `__clist_tmp:w` contains `\s__clist_mark`: in that case, `__clist_remove_all:w` removes the second `\s__clist_mark` (inserted by `__clist_tmp:w`), and lets `__clist_use_none_delimit_by_s_stop:w` act.

No brace is lost because items are always grabbed with a leading comma. The result of the first assignment has an extra leading comma, which we remove in a second assignment. Two exceptions: if the clist lost all of its elements, the result is empty, and we shouldn’t remove anything; if the clist started up empty, the first step happens to turn it into a single comma, and the second step removes it.

```

10068 \cs_new_protected:Npn \clist_remove_all:Nn
10069   { \__clist_remove_all:NNNn \clist_set_from_seq:NN \__kernel_tl_set:Nx }
10070 \cs_new_protected:Npn \clist_gremove_all:Nn
10071   { \__clist_remove_all:NNNn \clist_gset_from_seq:NN \__kernel_tl_gset:Nx }
10072 \cs_new_protected:Npn \__clist_remove_all:NNNn #1#2#3#4
10073   {
10074     \__clist_if_wrap:nTF {#4}
10075     {
10076       \seq_set_from_clist:NN \l__clist_internal_remove_seq #3
10077       \seq_remove_all:Nn \l__clist_internal_remove_seq {#4}
10078       #1 #3 \l__clist_internal_remove_seq
10079     }
10080     {
10081       \cs_set:Npn \__clist_tmp:w ##1 , #4 ,
10082       {
10083         ##1
10084         , \s__clist_mark , \__clist_use_none_delimit_by_s_stop:w ,
10085         \__clist_remove_all:
10086       }
10087       #2 #3
10088       {
10089         \exp_after:wN \__clist_remove_all:
10090         #3 , \s__clist_mark , #4 , \s__clist_stop
10091       }
10092       \clist_if_empty:NF #3
10093       {
10094         #2 #3
10095         {
10096           \exp_args:No \exp_not:o

```

```

10097         { \exp_after:wN \use_none:n #3 }
10098     }
10099 }
10100 }
10101 }
10102 \cs_new:Npn \__clist_remove_all:
10103 { \exp_after:wN \__clist_remove_all:w \__clist_tmp:w , }
10104 \cs_new:Npn \__clist_remove_all:w #1 , \s__clist_mark , #2 , { \exp_not:n {#1} }
10105 \cs_generate_variant:Nn \clist_remove_all:Nn { c }
10106 \cs_generate_variant:Nn \clist_gremove_all:Nn { c }

```

(End definition for `\clist_remove_all:Nn` and others. These functions are documented on page 122.)

`\clist_reverse:N` Use `\clist_reverse:n` in an x-expanding assignment. The extra work that `\clist_reverse:n` does to preserve braces and spaces would not be needed for the well-controlled case of N-type comma lists, but the slow-down is not too bad.

`\clist_reverse:c`
`\clist_greverse:N`
`\clist_greverse:c`

```

10107 \cs_new_protected:Npn \clist_reverse:N #1
10108 { \__kernel_tl_set:Nx #1 { \exp_args:No \clist_reverse:n {#1} } }
10109 \cs_new_protected:Npn \clist_greverse:N #1
10110 { \__kernel_tl_gset:Nx #1 { \exp_args:No \clist_reverse:n {#1} } }
10111 \cs_generate_variant:Nn \clist_reverse:N { c }
10112 \cs_generate_variant:Nn \clist_greverse:N { c }

```

(End definition for `\clist_reverse:N` and `\clist_greverse:N`. These functions are documented on page 122.)

`\clist_reverse:n` The reversed token list is built one item at a time, and stored between `\s__clist_stop` and `\s__clist_mark`, in the form of ? followed by zero or more instances of “`<item>`,”. We start from a comma list “`<item1>, ..., <itemn>`”. During the loop, the auxiliary `__clist_reverse:wwNww` receives “`?<itemi>`” as #1, “`<itemi+1>, ..., <itemn>`” as #2, `__clist_reverse:wwNww` as #3, what remains until `\s__clist_stop` as #4, and “`<itemi-1>, ..., <item1>`,” as #5. The auxiliary moves #1 just before #5, with a comma, and calls itself (#3). After the last item is moved, `__clist_reverse:wwNww` receives “`\s__clist_mark __clist_reverse:wwNww !`” as its argument #1, thus `__clist_reverse_end:ww` as its argument #3. This second auxiliary cleans up until the marker !, removes the trailing comma (introduced when the first item was moved after `\s__clist_stop`), and leaves its argument #1 within `\exp_not:n`. There is also a need to remove a leading comma, hence `\exp_not:o` and `\use_none:n`.

```

10113 \cs_new:Npn \clist_reverse:n #1
10114 {
10115     \__clist_reverse:wwNww ? #1 ,
10116     \s__clist_mark \__clist_reverse:wwNww ! ,
10117     \s__clist_mark \__clist_reverse_end:ww
10118     \s__clist_stop ? \s__clist_mark
10119 }
10120 \cs_new:Npn \__clist_reverse:wwNww
10121 #1 , #2 \s__clist_mark #3 #4 \s__clist_stop ? #5 \s__clist_mark
10122 { #3 ? #2 \s__clist_mark #3 #4 \s__clist_stop #1 , #5 \s__clist_mark }
10123 \cs_new:Npn \__clist_reverse_end:ww #1 ! #2 , \s__clist_mark
10124 { \exp_not:o { \use_none:n #2 } }

```

(End definition for `\clist_reverse:n`, `__clist_reverse:wwNww`, and `__clist_reverse_end:ww`. This function is documented on page 122.)

`\clist_sort:Nn` Implemented in `l3sort`.

`\clist_sort:cn`

`\clist_gsort:Nn` (End definition for `\clist_sort:Nn` and `\clist_gsort:Nn`. These functions are documented on page 122.)

`\clist_gsort:cn`

16.6 Comma list conditionals

`\clist_if_empty_p:N` Simple copies from the token list variable material.

`\clist_if_empty_p:c` 10125 `\prg_new_eq_conditional:Nn \clist_if_empty:N \tl_if_empty:N`

`\clist_if_empty:NTF` 10126 `{ p , T , F , TF }`

`\clist_if_empty:cTF` 10127 `\prg_new_eq_conditional:Nn \clist_if_empty:c \tl_if_empty:c`

10128 `{ p , T , F , TF }`

(End definition for `\clist_if_empty:N`. This function is documented on page 123.)

`\clist_if_empty_p:n` As usual, we insert a token (here ?) before grabbing any argument: this avoids losing braces. The argument of `\tl_if_empty:oTF` is empty if #1 is ? followed by blank spaces (besides, this particular variant of the emptiness test is optimized). If the item of the comma list is blank, grab the next one. As soon as one item is non-blank, exit: the second auxiliary grabs `\prg_return_false:` as #2, unless every item in the comma list was blank and the loop actually got broken by the trailing `\s__clist_mark \prg_return_false:` item.

`\clist_if_empty:nTF`

`__clist_if_empty_n:w`

`__clist_if_empty_n:wNw`

```
10129 \prg_new_conditional:Npnn \clist_if_empty:n #1 { p , T , F , TF }
10130 {
10131   \__clist_if_empty_n:w ? #1
10132   , \s__clist_mark \prg_return_false:
10133   , \s__clist_mark \prg_return_true:
10134   \s__clist_stop
10135 }
10136 \cs_new:Npn \__clist_if_empty_n:w #1 ,
10137 {
10138   \tl_if_empty:oTF { \use_none:nn #1 ? }
10139   { \__clist_if_empty_n:w ? }
10140   { \__clist_if_empty_n:wNw }
10141 }
10142 \cs_new:Npn \__clist_if_empty_n:wNw #1 \s__clist_mark #2#3 \s__clist_stop {#2}
```

(End definition for `\clist_if_empty:nTF`, `__clist_if_empty_n:w`, and `__clist_if_empty_n:wNw`. This function is documented on page 123.)

`\clist_if_in:NnTF` For “safe” items, we simply surround the comma list, and the item, with commas, then use the same code as for `\tl_if_in:Nn`. For “unsafe” items we follow the same route as `\seq_if_in:Nn`, mapping through the list a comparison function. If found, return `true` and remove `\prg_return_false:`.

`\clist_if_in:NVTF` 10143 `\prg_new_protected_conditional:Npnn \clist_if_in:Nn #1#2 { T , F , TF }`

`\clist_if_in:NoTF` 10144 `{`

`\clist_if_in:cnTF` 10145 `\exp_args:No __clist_if_in_return:nnN #1 {#2} #1`

`\clist_if_in:cVTF` 10146 `}`

`\clist_if_in:coTF` 10147 `\prg_new_protected_conditional:Npnn \clist_if_in:nn #1#2 { T , F , TF }`

`\clist_if_in:nnTF` 10148 `{`

`__clist_if_in_return:nnN` 10149 `\clist_set:Nn \l__clist_internal_clist {#1}`

10150 `\exp_args:No __clist_if_in_return:nnN \l__clist_internal_clist {#2}`

10151 `\l__clist_internal_clist`

10152 `}`

```

10153 \cs_new_protected:Npn \__clist_if_in_return:nnN #1#2#3
10154 {
10155     \__clist_if_wrap:nTF {#2}
10156     {
10157         \cs_set:Npx \__clist_tmp:w ##1
10158         {
10159             \exp_not:N \tl_if_eq:nnT {##1}
10160             \exp_not:n
10161             {
10162                 {#2}
10163                 { \clist_map_break:n { \prg_return_true: \use_none:n } }
10164             }
10165         }
10166         \clist_map_function:NN #3 \__clist_tmp:w
10167         \prg_return_false:
10168     }
10169     {
10170         \cs_set:Npn \__clist_tmp:w ##1 ,#2, { }
10171         \tl_if_empty:oTF
10172         { \__clist_tmp:w ,#1, {} {} ,#2, }
10173         { \prg_return_false: } { \prg_return_true: }
10174     }
10175 }
10176 \prg_generate_conditional_variant:Nnn \clist_if_in:Nn
10177 { NV , No , c , cV , co } { T , F , TF }
10178 \prg_generate_conditional_variant:Nnn \clist_if_in:nn
10179 { nV , no } { T , F , TF }

```

(End definition for `\clist_if_in:NnTF`, `\clist_if_in:nnTF`, and `__clist_if_in_return:nnN`. These functions are documented on page 123.)

16.7 Mapping to comma lists

`\clist_map_function:NN`
`\clist_map_function:cN`
`__clist_map_function:Nw`

If the variable is empty, the mapping is skipped (otherwise, that comma-list would be seen as consisting of one empty item). Then loop over the comma-list, grabbing one comma-delimited item at a time. The end is marked by `\q__clist_recursion_tail`. The auxiliary function `__clist_map_function:Nw` is also used in `\clist_map_inline:Nn`.

```

10180 \cs_new:Npn \clist_map_function:NN #1#2
10181 {
10182     \clist_if_empty:NF #1
10183     {
10184         \exp_last_unbraced:NNo \__clist_map_function:Nw #2 #1
10185         , \q__clist_recursion_tail ,
10186         \prg_break_point:Nn \clist_map_break: { }
10187     }
10188 }
10189 \cs_new:Npn \__clist_map_function:Nw #1#2 ,
10190 {
10191     \__clist_if_recursion_tail_break:nN {#2} \clist_map_break:
10192     #1 {#2}
10193     \__clist_map_function:Nw #1
10194 }
10195 \cs_generate_variant:Nn \clist_map_function:NN { c }

```

(End definition for `\clist_map_function:Nn` and `__clist_map_function:Nw`. This function is documented on page 123.)

```

\clist_map_function:nN The n-type mapping function is a bit more awkward, since spaces must be trimmed from
__clist_map_function_n:Nn each item. Space trimming is again based on __clist_trim_next:w. The auxiliary
__clist_map_unbrace:Nw __clist_map_function_n:Nn receives as arguments the function, and the next non-
empty item (after space trimming but before brace removal). One level of braces is
removed by __clist_map_unbrace:Nw.

10196 \cs_new:Npn \clist_map_function:nN #1#2
10197 {
10198     \exp_after:wN __clist_map_function_n:Nn \exp_after:wN #2
10199     \exp:w __clist_trim_next:w \prg_do_nothing: #1 , \q__clist_recursion_tail ,
10200     \prg_break_point:Nn \clist_map_break: { }
10201 }
10202 \cs_new:Npn __clist_map_function_n:Nn #1 #2
10203 {
10204     __clist_if_recursion_tail_break:nN {#2} \clist_map_break:
10205     __clist_map_unbrace:Nw #1 #2,
10206     \exp_after:wN __clist_map_function_n:Nn \exp_after:wN #1
10207     \exp:w __clist_trim_next:w \prg_do_nothing:
10208 }
10209 \cs_new:Npn __clist_map_unbrace:Nw #1 #2, { #1 {#2} }

```

(End definition for `\clist_map_function:nN`, `__clist_map_function_n:Nn`, and `__clist_map_unbrace:Nw`. This function is documented on page 123.)

`\clist_map_inline:Nn` Inline mapping is done by creating a suitable function “on the fly”: this is done globally
`\clist_map_inline:cn` to avoid any issues with \TeX ’s groups. We use a different function for each level of
`\clist_map_inline:nn` nesting.

Since the mapping is non-expandable, we can perform the space-trimming needed by the `n` version simply by storing the comma-list in a variable. We don’t need a different comma-list for each nesting level: the comma-list is expanded before the mapping starts.

```

10210 \cs_new_protected:Npn \clist_map_inline:Nn #1#2
10211 {
10212     \clist_if_empty:NF #1
10213     {
10214         \int_gincr:N \g__kernel_prg_map_int
10215         \cs_gset_protected:cpn
10216         { __clist_map_ \int_use:N \g__kernel_prg_map_int :w } ##1 {#2}
10217         \exp_last_unbraced:Nco __clist_map_function:Nw
10218         { __clist_map_ \int_use:N \g__kernel_prg_map_int :w }
10219         #1 , \q__clist_recursion_tail ,
10220         \prg_break_point:Nn \clist_map_break:
10221         { \int_gdecr:N \g__kernel_prg_map_int }
10222     }
10223 }
10224 \cs_new_protected:Npn \clist_map_inline:nn #1
10225 {
10226     \clist_set:Nn \l__clist_internal_clist {#1}
10227     \clist_map_inline:Nn \l__clist_internal_clist
10228 }
10229 \cs_generate_variant:Nn \clist_map_inline:Nn { c }

```

(End definition for `\clist_map_inline:Nn` and `\clist_map_inline:nn`. These functions are documented on page 124.)

`\clist_map_variable:NNn` As for other comma-list mappings, filter out the case of an empty list. Same approach as
`\clist_map_variable:cNn` `\clist_map_function:Nn`, additionally we store each item in the given variable. As for
`\clist_map_variable:nNn` inline mappings, space trimming for the `n` variant is done by storing the comma list in
`__clist_map_variable:Nnw` a variable. The quark test is done before assigning the item to the variable: this avoids
storing a quark which the user wouldn't expect. The strange `\use:n` avoids unlikely
problems when `#2` would contain `\q__clist_recursion_stop`.

```

10230 \cs_new_protected:Npn \clist_map_variable:NNn #1#2#3
10231 {
10232   \clist_if_empty:NF #1
10233   {
10234     \exp_args:Nno \use:n
10235     { \__clist_map_variable:Nnw #2 {#3} }
10236     #1
10237     , \q__clist_recursion_tail , \q__clist_recursion_stop
10238     \prg_break_point:Nn \clist_map_break: { }
10239   }
10240 }
10241 \cs_new_protected:Npn \clist_map_variable:nNn #1
10242 {
10243   \clist_set:Nn \l__clist_internal_clist {#1}
10244   \clist_map_variable:NNn \l__clist_internal_clist
10245 }
10246 \cs_new_protected:Npn \__clist_map_variable:Nnw #1#2#3,
10247 {
10248   \__clist_if_recursion_tail_stop:n {#3}
10249   \tl_set:Nn #1 {#3}
10250   \use:n {#2}
10251   \__clist_map_variable:Nnw #1 {#2}
10252 }
10253 \cs_generate_variant:Nn \clist_map_variable:NNn { c }

```

(End definition for `\clist_map_variable:NNn`, `\clist_map_variable:nNn`, and `__clist_map_variable:Nnw`.
These functions are documented on page 124.)

`\clist_map_break:` The break statements use the general `\prg_map_break:Nn` mechanism.
`\clist_map_break:n`

```

10254 \cs_new:Npn \clist_map_break:
10255 { \prg_map_break:Nn \clist_map_break: { } }
10256 \cs_new:Npn \clist_map_break:n
10257 { \prg_map_break:Nn \clist_map_break: }

```

(End definition for `\clist_map_break:` and `\clist_map_break:n`. These functions are documented on
page 124.)

`\clist_count:N` Counting the items in a comma list is done using the same approach as for other token
`\clist_count:c` count functions: turn each entry into a `+1` then use integer evaluation to actually do the
`\clist_count:n` mathematics. In the case of an `n`-type comma-list, we could of course use `\clist_map_-`
`__clist_count:n` `function:nN`, but that is very slow, because it carefully removes spaces. Instead, we loop
`__clist_count:w` manually, and skip blank items (but not `{}`, hence the extra spaces).

```

10258 \cs_new:Npn \clist_count:N #1
10259 {
10260   \int_eval:n
10261   {
10262     0
10263     \clist_map_function:NN #1 \__clist_count:n

```



```

10264     }
10265   }
10266   \cs_generate_variant:Nn \clist_count:N { c }
10267   \cs_new:Npx \clist_count:n #1
10268   {
10269     \exp_not:N \int_eval:n
10270     {
10271       0
10272       \exp_not:N \__clist_count:w \c_space_tl
10273       #1 \exp_not:n { , \q__clist_recursion_tail , \q__clist_recursion_stop }
10274     }
10275   }
10276   \cs_new:Npn \__clist_count:n #1 { + 1 }
10277   \cs_new:Npx \__clist_count:w #1 ,
10278   {
10279     \exp_not:n { \exp_args:Nf \__clist_if_recursion_tail_stop:n } {#1}
10280     \exp_not:N \tl_if_blank:nF {#1} { + 1 }
10281     \exp_not:N \__clist_count:w \c_space_tl
10282   }

```

(End definition for `\clist_count:N` and others. These functions are documented on page 125.)

16.8 Using comma lists

`\clist_use:Nnnn` First check that the variable exists. Then count the items in the comma list. If it has none, output nothing. If it has one item, output that item, brace stripped (note that space-trimming has already been done when the comma list was assigned). If it has two, place the *<separator between two>* in the middle.

`\clist_use:cnnn` Otherwise, `__clist_use:nwwwnwn` takes the following arguments; 1: a *<separator>*, 2, 3, 4: three items from the comma list (or quarks), 5: the rest of the comma list, 6: a *<continuation>* function (`use_ii` or `use_iii` with its *<separator>* argument), 7: junk, and 8: the temporary result, which is built in a brace group following `\q__clist_stop`.

`__clist_use:wwn` The *<separator>* and the first of the three items are placed in the result, then we use the *<continuation>*, placing the remaining two items after it. When we begin this loop, the three items really belong to the comma list, the first `\q__clist_mark` is taken as a delimiter to the `use_ii` function, and the continuation is `use_ii` itself. When we reach the last two items of the original token list, `\q__clist_mark` is taken as a third item, and now the second `\q__clist_mark` serves as a delimiter to `use_ii`, switching to the other *<continuation>*, `use_iii`, which uses the *<separator between final two>*.

`__clist_use:nwwwn`

`\clist_use:Nn`

`\clist_use:cn`

```

10283   \cs_new:Npn \clist_use:Nnnn #1#2#3#4
10284   {
10285     \clist_if_exist:NTF #1
10286     {
10287       \int_case:nnF { \clist_count:N #1 }
10288       {
10289         { 0 } { }
10290         { 1 } { \exp_after:wN \__clist_use:wwn #1 , , { } }
10291         { 2 } { \exp_after:wN \__clist_use:wwn #1 , {#2} }
10292       }
10293       {
10294         \exp_after:wN \__clist_use:nwwwnwn
10295         \exp_after:wN { \exp_after:wN } #1 ,
10296         \s__clist_mark , { \__clist_use:nwwwnwn {#3} }

```

```

10297         \s__clist_mark , { \__clist_use:nwn {#4} }
10298         \s__clist_stop { }
10299     }
10300 }
10301 {
10302     \__kernel_msg_expandable_error:nnn
10303     { kernel } { bad-variable } {#1}
10304 }
10305 }
10306 \cs_generate_variant:Nn \clist_use:Nnnn { c }
10307 \cs_new:Npn \__clist_use:wn #1 , #2 , #3 { \exp_not:n { #1 #3 #2 } }
10308 \cs_new:Npn \__clist_use:nwnnnwn
10309     #1#2 , #3 , #4 , #5 \s__clist_mark , #6#7 \s__clist_stop #8
10310     { #6 {#3} , {#4} , #5 \s__clist_mark , {#6} #7 \s__clist_stop { #8 #1 #2 } }
10311 \cs_new:Npn \__clist_use:nwn #1#2 , #3 \s__clist_stop #4
10312     { \exp_not:n { #4 #1 #2 } }
10313 \cs_new:Npn \clist_use:Nn #1#2
10314     { \clist_use:Nnnn #1 {#2} {#2} {#2} }
10315 \cs_generate_variant:Nn \clist_use:Nn { c }

```

(End definition for `\clist_use:Nnnn` and others. These functions are documented on page 125.)

16.9 Using a single item

`\clist_item:Nn` To avoid needing to test the end of the list at each step, we first compute the $\langle length \rangle$ of the list. If the item number is 0, less than $-\langle length \rangle$, or more than $\langle length \rangle$, the result is empty. If it is negative, but not less than $-\langle length \rangle$, add $\langle length \rangle + 1$ to the item number before performing the loop. The loop itself is very simple, return the item if the counter reached 1, otherwise, decrease the counter and repeat.

`\clist_item:cn`

`__clist_item:nnnN`

`__clist_item:ffoN`

`__clist_item:ffnN`

`__clist_item_N_loop:nw`

```

10316 \cs_new:Npn \clist_item:Nn #1#2
10317 {
10318     \__clist_item:ffoN
10319     { \clist_count:N #1 }
10320     { \int_eval:n {#2} }
10321     #1
10322     \__clist_item_N_loop:nw
10323 }
10324 \cs_new:Npn \__clist_item:nnnN #1#2#3#4
10325 {
10326     \int_compare:nNnTF {#2} < 0
10327     {
10328         \int_compare:nNnTF {#2} < { - #1 }
10329         { \__clist_use_none_delimit_by_s_stop:w }
10330         { \exp_args:Nf #4 { \int_eval:n { #2 + 1 + #1 } } }
10331     }
10332     {
10333         \int_compare:nNnTF {#2} > {#1}
10334         { \__clist_use_none_delimit_by_s_stop:w }
10335         { #4 {#2} }
10336     }
10337     { } , #3 , \s__clist_stop
10338 }
10339 \cs_generate_variant:Nn \__clist_item:nnnN { ffo, ff }
10340 \cs_new:Npn \__clist_item_N_loop:nw #1 #2,

```

```

10341 {
10342   \int_compare:nNnTF {#1} = 0
10343     { \__clist_use_i_delimit_by_s_stop:nw { \exp_not:n {#2} } }
10344     { \exp_args:Nf \__clist_item_N_loop:nw { \int_eval:n { #1 - 1 } } }
10345   }
10346 \cs_generate_variant:Nn \clist_item:Nn { c }

```

(End definition for `\clist_item:Nn`, `__clist_item:nnnN`, and `__clist_item_N_loop:nw`. This function is documented on page 127.)

`\clist_item:nn` This starts in the same way as `\clist_item:Nn` by counting the items of the comma list. The final item should be space-trimmed before being brace-stripped, hence we insert a couple of odd-looking `\prg_do_nothing:` to avoid losing braces. Blank items are ignored.

```

\__clist_item_n:nw
\__clist_item_n_loop:nw
\__clist_item_n_end:n
\__clist_item_n_strip:n
\__clist_item_n_strip:w
10347 \cs_new:Npn \clist_item:nn #1#2
10348 {
10349   \__clist_item:ffnN
10350   { \clist_count:n {#1} }
10351   { \int_eval:n {#2} }
10352   {#1}
10353   \__clist_item_n:nw
10354 }
10355 \cs_new:Npn \__clist_item_n:nw #1
10356 { \__clist_item_n_loop:nw {#1} \prg_do_nothing: }
10357 \cs_new:Npn \__clist_item_n_loop:nw #1 #2,
10358 {
10359   \exp_args:No \tl_if_blank:nTF {#2}
10360   { \__clist_item_n_loop:nw {#1} \prg_do_nothing: }
10361   {
10362     \int_compare:nNnTF {#1} = 0
10363       { \exp_args:No \__clist_item_n_end:n {#2} }
10364       {
10365         \exp_args:Nf \__clist_item_n_loop:nw
10366         { \int_eval:n { #1 - 1 } }
10367         \prg_do_nothing:
10368       }
10369   }
10370 }
10371 \cs_new:Npn \__clist_item_n_end:n #1 #2 \s_clist_stop
10372 { \tl_trim_spaces_apply:nN {#1} \__clist_item_n_strip:n }
10373 \cs_new:Npn \__clist_item_n_strip:n #1 { \__clist_item_n_strip:w #1 , }
10374 \cs_new:Npn \__clist_item_n_strip:w #1 , { \exp_not:n {#1} }

```

(End definition for `\clist_item:nn` and others. This function is documented on page 127.)

`\clist_rand_item:n` The N-type function is not implemented through the n-type function for efficiency: for instance comma-list variables do not require space-trimming of their items. Even testing for emptiness of an n-type comma-list is slow, so we count items first and use that both for the emptiness test and the pseudo-random integer. Importantly, `\clist_item:Nn` and `\clist_item:nn` only evaluate their argument once.

```

10375 \cs_new:Npn \clist_rand_item:n #1
10376 { \exp_args:Nf \__clist_rand_item:nn { \clist_count:n {#1} } {#1} }
10377 \cs_new:Npn \__clist_rand_item:nn #1#2
10378 {
10379   \int_compare:nNnF {#1} = 0

```

```

10380     { \clist_item:nn {#2} { \int_rand:nn { 1 } {#1} } }
10381   }
10382   \cs_new:Npn \clist_rand_item:N #1
10383   {
10384     \clist_if_empty:NF #1
10385     { \clist_item:Nn #1 { \int_rand:nn { 1 } { \clist_count:N #1 } } }
10386   }
10387   \cs_generate_variant:Nn \clist_rand_item:N { c }

```

(End definition for `\clist_rand_item:n`, `\clist_rand_item:N`, and `__clist_rand_item:nn`. These functions are documented on page 127.)

16.10 Viewing comma lists

`\clist_show:N` Apply the general `__kernel_chk_defined:NT` and `\msg_show:nnnnnn`.

```

\clist_show:c 10388 \cs_new_protected:Npn \clist_show:N { \__clist_show:NN \msg_show:nnxxxx }
\clist_log:N   10389 \cs_generate_variant:Nn \clist_show:N { c }
\clist_log:c   10390 \cs_new_protected:Npn \clist_log:N { \__clist_show:NN \msg_log:nnxxxx }
\__clist_show:NN 10391 \cs_generate_variant:Nn \clist_log:N { c }
                10392 \cs_new_protected:Npn \__clist_show:NN #1#2
                10393 {
                10394   \__kernel_chk_defined:NT #2
                10395   {
                10396     #1 { LaTeX/kernel } { show-clist }
                10397     { \token_to_str:N #2 }
                10398     { \clist_map_function:NN #2 \msg_show_item:n }
                10399     { } { }
                10400   }
                10401 }

```

(End definition for `\clist_show:N`, `\clist_log:N`, and `__clist_show:NN`. These functions are documented on page 127.)

`\clist_show:n` A variant of the above: no existence check, empty first argument for the message.

```

\clist_log:n   10402 \cs_new_protected:Npn \clist_show:n { \__clist_show:NN \msg_show:nnxxxx }
\__clist_show:NN 10403 \cs_new_protected:Npn \clist_log:n { \__clist_show:NN \msg_log:nnxxxx }
                10404 \cs_new_protected:Npn \__clist_show:NN #1#2
                10405 {
                10406   #1 { LaTeX/kernel } { show-clist }
                10407   { } { \clist_map_function:nN {#2} \msg_show_item:n } { } { }
                10408 }

```

(End definition for `\clist_show:n`, `\clist_log:n`, and `__clist_show:NN`. These functions are documented on page 128.)

16.11 Scratch comma lists

`\l_tmpa_clist` Temporary comma list variables.

```

\l_tmpb_clist 10409 \clist_new:N \l_tmpa_clist
\g_tmpa_clist 10410 \clist_new:N \l_tmpb_clist
\g_tmpb_clist 10411 \clist_new:N \g_tmpa_clist
                10412 \clist_new:N \g_tmpb_clist

```

(End definition for `\l_tmpa_clist` and others. These variables are documented on page 128.)

```

10413 </package>

```

17 l3token implementation

```
10414 \*package>
10415 \*tex>
10416 \@@=char>
```

17.1 Internal auxiliaries

`\s__char_stop` Internal scan mark.

```
10417 \scan_new:N \s__char_stop
```

(End definition for `\s__char_stop`.)

`\q__char_no_value` Internal recursion quarks.

```
10418 \quark_new:N \q__char_no_value
```

(End definition for `\q__char_no_value`.)

`_char_quark_if_no_value_p:N` Functions to query recursion quarks.

```
\_char_quark_if_no_value:NTF 10419 \_kernel_quark_new_conditional:Nn \_char_quark_if_no_value:N { TF }
```

(End definition for `_char_quark_if_no_value:N`.)

17.2 Manipulating and interrogating character tokens

`\char_set_catcode:nn` Simple wrappers around the primitives.

```
\char_value_catcode:n 10420 \cs_new_protected:Npn \char_set_catcode:nn #1#2
\char_show_value_catcode:n 10421 { \tex_catcode:D \int_eval:n {#1} = \int_eval:n {#2} \exp_stop_f: }
10422 \cs_new:Npn \char_value_catcode:n #1
10423 { \tex_the:D \tex_catcode:D \int_eval:n {#1} \exp_stop_f: }
10424 \cs_new_protected:Npn \char_show_value_catcode:n #1
10425 { \exp_args:Nf \tl_show:n { \char_value_catcode:n {#1} } }
```

(End definition for `\char_set_catcode:nn`, `\char_value_catcode:n`, and `\char_show_value_catcode:n`. These functions are documented on page 132.)

`\char_set_catcode_escape:N`

`\char_set_catcode_group_begin:N`

`\char_set_catcode_group_end:N`

`\char_set_catcode_math_toggle:N`

`\char_set_catcode_alignment:N`

`\char_set_catcode_end_line:N`

`\char_set_catcode_parameter:N`

`\char_set_catcode_math_superscript:N`

`\char_set_catcode_math_subscript:N`

`\char_set_catcode_ignore:N`

`\char_set_catcode_space:N`

`\char_set_catcode_letter:N`

`\char_set_catcode_other:N`

`\char_set_catcode_active:N`

`\char_set_catcode_comment:N`

`\char_set_catcode_invalid:N`

```
10426 \cs_new_protected:Npn \char_set_catcode_escape:N #1
10427 { \char_set_catcode:nn { '#1 } { 0 } }
10428 \cs_new_protected:Npn \char_set_catcode_group_begin:N #1
10429 { \char_set_catcode:nn { '#1 } { 1 } }
10430 \cs_new_protected:Npn \char_set_catcode_group_end:N #1
10431 { \char_set_catcode:nn { '#1 } { 2 } }
10432 \cs_new_protected:Npn \char_set_catcode_math_toggle:N #1
10433 { \char_set_catcode:nn { '#1 } { 3 } }
10434 \cs_new_protected:Npn \char_set_catcode_alignment:N #1
10435 { \char_set_catcode:nn { '#1 } { 4 } }
10436 \cs_new_protected:Npn \char_set_catcode_end_line:N #1
10437 { \char_set_catcode:nn { '#1 } { 5 } }
10438 \cs_new_protected:Npn \char_set_catcode_parameter:N #1
10439 { \char_set_catcode:nn { '#1 } { 6 } }
10440 \cs_new_protected:Npn \char_set_catcode_math_superscript:N #1
10441 { \char_set_catcode:nn { '#1 } { 7 } }
10442 \cs_new_protected:Npn \char_set_catcode_math_subscript:N #1
```

```

10443 { \char_set_catcode:nn { '#1 } { 8 } }
10444 \cs_new_protected:Npn \char_set_catcode_ignore:N #1
10445 { \char_set_catcode:nn { '#1 } { 9 } }
10446 \cs_new_protected:Npn \char_set_catcode_space:N #1
10447 { \char_set_catcode:nn { '#1 } { 10 } }
10448 \cs_new_protected:Npn \char_set_catcode_letter:N #1
10449 { \char_set_catcode:nn { '#1 } { 11 } }
10450 \cs_new_protected:Npn \char_set_catcode_other:N #1
10451 { \char_set_catcode:nn { '#1 } { 12 } }
10452 \cs_new_protected:Npn \char_set_catcode_active:N #1
10453 { \char_set_catcode:nn { '#1 } { 13 } }
10454 \cs_new_protected:Npn \char_set_catcode_comment:N #1
10455 { \char_set_catcode:nn { '#1 } { 14 } }
10456 \cs_new_protected:Npn \char_set_catcode_invalid:N #1
10457 { \char_set_catcode:nn { '#1 } { 15 } }

```

(End definition for `\char_set_catcode_escape:N` and others. These functions are documented on page 131.)

```

\char_set_catcode_escape:n
\char_set_catcode_group_begin:n
\char_set_catcode_group_end:n
\char_set_catcode_math_toggle:n
\char_set_catcode_alignment:n
\char_set_catcode_end_line:n
\char_set_catcode_parameter:n
\char_set_catcode_math_superscript:n
\char_set_catcode_math_subscript:n
\char_set_catcode_ignore:n
\char_set_catcode_space:n
\char_set_catcode_letter:n
\char_set_catcode_other:n
\char_set_catcode_active:n
\char_set_catcode_comment:n
\char_set_catcode_invalid:n

10458 \cs_new_protected:Npn \char_set_catcode_escape:n #1
10459 { \char_set_catcode:nn {#1} { 0 } }
10460 \cs_new_protected:Npn \char_set_catcode_group_begin:n #1
10461 { \char_set_catcode:nn {#1} { 1 } }
10462 \cs_new_protected:Npn \char_set_catcode_group_end:n #1
10463 { \char_set_catcode:nn {#1} { 2 } }
10464 \cs_new_protected:Npn \char_set_catcode_math_toggle:n #1
10465 { \char_set_catcode:nn {#1} { 3 } }
10466 \cs_new_protected:Npn \char_set_catcode_alignment:n #1
10467 { \char_set_catcode:nn {#1} { 4 } }
10468 \cs_new_protected:Npn \char_set_catcode_end_line:n #1
10469 { \char_set_catcode:nn {#1} { 5 } }
10470 \cs_new_protected:Npn \char_set_catcode_parameter:n #1
10471 { \char_set_catcode:nn {#1} { 6 } }
10472 \cs_new_protected:Npn \char_set_catcode_math_superscript:n #1
10473 { \char_set_catcode:nn {#1} { 7 } }
10474 \cs_new_protected:Npn \char_set_catcode_math_subscript:n #1
10475 { \char_set_catcode:nn {#1} { 8 } }
10476 \cs_new_protected:Npn \char_set_catcode_ignore:n #1
10477 { \char_set_catcode:nn {#1} { 9 } }
10478 \cs_new_protected:Npn \char_set_catcode_space:n #1
10479 { \char_set_catcode:nn {#1} { 10 } }
10480 \cs_new_protected:Npn \char_set_catcode_letter:n #1
10481 { \char_set_catcode:nn {#1} { 11 } }
10482 \cs_new_protected:Npn \char_set_catcode_other:n #1
10483 { \char_set_catcode:nn {#1} { 12 } }
10484 \cs_new_protected:Npn \char_set_catcode_active:n #1
10485 { \char_set_catcode:nn {#1} { 13 } }
10486 \cs_new_protected:Npn \char_set_catcode_comment:n #1
10487 { \char_set_catcode:nn {#1} { 14 } }
10488 \cs_new_protected:Npn \char_set_catcode_invalid:n #1
10489 { \char_set_catcode:nn {#1} { 15 } }

```

(End definition for `\char_set_catcode_escape:n` and others. These functions are documented on page 131.)

```

\char_set_mathcode:nn
\char_value_mathcode:n
\char_show_value_mathcode:n
\char_set_lccode:nn
\char_value_lccode:n
\char_show_value_lccode:n
\char_set_uccode:nn
\char_value_uccode:n
\char_show_value_uccode:n
\char_set_sfcode:nn
\char_value_sfcode:n
\char_show_value_sfcode:n

```

Pretty repetitive, but necessary!

```

10490 \cs_new_protected:Npn \char_set_mathcode:nn #1#2
10491 { \tex_mathcode:D \int_eval:n {#1} = \int_eval:n {#2} \exp_stop_f: }
10492 \cs_new:Npn \char_value_mathcode:n #1
10493 { \tex_the:D \tex_mathcode:D \int_eval:n {#1} \exp_stop_f: }
10494 \cs_new_protected:Npn \char_show_value_mathcode:n #1
10495 { \exp_args:Nf \tl_show:n { \char_value_mathcode:n {#1} } }
10496 \cs_new_protected:Npn \char_set_lccode:nn #1#2
10497 { \tex_lccode:D \int_eval:n {#1} = \int_eval:n {#2} \exp_stop_f: }
10498 \cs_new:Npn \char_value_lccode:n #1
10499 { \tex_the:D \tex_lccode:D \int_eval:n {#1} \exp_stop_f: }
10500 \cs_new_protected:Npn \char_show_value_lccode:n #1
10501 { \exp_args:Nf \tl_show:n { \char_value_lccode:n {#1} } }
10502 \cs_new_protected:Npn \char_set_uccode:nn #1#2
10503 { \tex_uccode:D \int_eval:n {#1} = \int_eval:n {#2} \exp_stop_f: }
10504 \cs_new:Npn \char_value_uccode:n #1
10505 { \tex_the:D \tex_uccode:D \int_eval:n {#1} \exp_stop_f: }
10506 \cs_new_protected:Npn \char_show_value_uccode:n #1
10507 { \exp_args:Nf \tl_show:n { \char_value_uccode:n {#1} } }
10508 \cs_new_protected:Npn \char_set_sfcode:nn #1#2
10509 { \tex_sfcode:D \int_eval:n {#1} = \int_eval:n {#2} \exp_stop_f: }
10510 \cs_new:Npn \char_value_sfcode:n #1
10511 { \tex_the:D \tex_sfcode:D \int_eval:n {#1} \exp_stop_f: }
10512 \cs_new_protected:Npn \char_show_value_sfcode:n #1
10513 { \exp_args:Nf \tl_show:n { \char_value_sfcode:n {#1} } }

```

(End definition for `\char_set_mathcode:nn` and others. These functions are documented on page 133.)

```

\l_char_active_seq
\l_char_special_seq

```

Two sequences for dealing with special characters. The first is characters which may be active, the second longer list is for “special” characters more generally. Both lists are escaped so that for example bulk code assignments can be carried out. In both cases, the order is by ASCII character code (as is done in for example `\ExplSyntaxOn`).

```

10514 \seq_new:N \l_char_special_seq
10515 \seq_set_split:Nnn \l_char_special_seq { }
10516 { \ \ " \# \$ \% \& \ \ ^ \_ \{ \} \~ }
10517 \seq_new:N \l_char_active_seq
10518 \seq_set_split:Nnn \l_char_active_seq { }
10519 { \ " \$ \& \ ^ \_ \~ }

```

(End definition for `\l_char_active_seq` and `\l_char_special_seq`. These variables are documented on page 133.)

17.3 Creating character tokens

```

\char_set_active_eq:NN
\char_set_active_eq:Nc
\char_gset_active_eq:NN
\char_gset_active_eq:Nc
\char_set_active_eq:nN
\char_set_active_eq:nc
\char_gset_active_eq:nN
\char_gset_active_eq:nc

```

Four simple functions with very similar definitions, so set up using an auxiliary. These are similar to LuaTeX’s `\letcharcode` primitive.

```

10520 \group_begin:
10521 \char_set_catcode_active:N \^^@
10522 \cs_set_protected:Npn \__char_tmp:nN #1#2
10523 {
10524 \cs_new_protected:cpn { #1 :nN } ##1
10525 {
10526 \group_begin:
10527 \char_set_lccode:nn { \^^@ } { ##1 }

```

```

10528         \tex_lowercase:D { \group_end: #2 ^^@ }
10529     }
10530     \cs_new_protected:cpx { #1 :NN } ##1
10531     { \exp_not:c { #1 : nN } { '##1 } }
10532 }
10533 \__char_tmp:nN { char_set_active_eq } \cs_set_eq:NN
10534 \__char_tmp:nN { char_gset_active_eq } \cs_gset_eq:NN
10535 \group_end:
10536 \cs_generate_variant:Nn \char_set_active_eq:NN { Nc }
10537 \cs_generate_variant:Nn \char_gset_active_eq:NN { Nc }
10538 \cs_generate_variant:Nn \char_set_active_eq:nN { nc }
10539 \cs_generate_variant:Nn \char_gset_active_eq:nN { nc }

```

(End definition for \char_set_active_eq:NN and others. These functions are documented on page 129.)

__char_int_to_roman:w For efficiency in 8-bit engines, we use the faster primitive approach to making roman numerals.

```

10540 \cs_new_eq:NN \__char_int_to_roman:w \tex_romannumeral:D

```

(End definition for __char_int_to_roman:w.)

\char_generate:nn The aim here is to generate characters of (broadly) arbitrary category code. Where possible, that is done using engine support (Xe_{La}TeX, Lua_{TeX}). There are though various issues which are covered below. At the interface layer, turn the two arguments into integers up-front so this is only done once.

```

\__char_generate_aux:nn
\__char_generate_aux:nnw
\__char_generate_auxii:nnw
  \l__char_tmp_tl
  \__char_generate_invalid_catcode:
10541 \cs_new:Npn \char_generate:nn #1#2
10542 {
10543   \exp:w \exp_after:wN \__char_generate_aux:w
10544   \int_value:w \int_eval:n {#1} \exp_after:wN ;
10545   \int_value:w \int_eval:n {#2} ;
10546 }

```

Before doing any actual conversion, first some special case filtering. Spaces are out here as Lua_{TeX} emulation only makes normal (charcode 32 spaces). However, ^^@ is filtered out separately as that can't be done with macro emulation either, so is flagged up separately. That done, hand off to the engine-dependent part.

```

10547 \cs_new:Npn \__char_generate_aux:w #1 ; #2 ;
10548 {
10549   \if_int_compare:w #2 = 10 \exp_stop_f:
10550   \if_int_compare:w #1 = 0 \exp_stop_f:
10551     \__kernel_msg_expandable_error:nn { kernel } { char-null-space }
10552   \else:
10553     \__kernel_msg_expandable_error:nn { kernel } { char-space }
10554   \fi:
10555   \else:
10556     \if_int_odd:w 0
10557       \if_int_compare:w #2 < 1 \exp_stop_f: 1 \fi:
10558       \if_int_compare:w #2 = 5 \exp_stop_f: 1 \fi:
10559       \if_int_compare:w #2 = 9 \exp_stop_f: 1 \fi:
10560       \if_int_compare:w #2 > 13 \exp_stop_f: 1 \fi: \exp_stop_f:
10561       \__kernel_msg_expandable_error:nn { kernel }
10562       { char-invalid-catcode }
10563     \else:
10564       \if_int_odd:w 0

```



```

10565         \if_int_compare:w #1 < 0 \exp_stop_f: 1 \fi:
10566         \if_int_compare:w #1 > \c_max_char_int 1 \fi: \exp_stop_f:
10567         \__kernel_msg_expandable_error:nn { kernel }
10568         { char-out-of-range }
10569     \else:
10570         \__char_generate_aux:nnw {#1} {#2}
10571     \fi:
10572 \fi:
10573 \fi:
10574 \exp_end:
10575 }
10576 \tl_new:N \l__char_tmp_tl

```

Engine-dependent definitions are now needed for the implementation. For LuaTeX and XeTeX there is engine-level support. They can do cases that macro emulation can't. All of those are filtered out here using a primitive-based boolean expression to avoid fixing the category code of the null character used in the false branch (for 8-bit engines). The final level is the basic definition at the engine level: the arguments here are integers so there is no need to worry about them too much. Older versions of XeTeX cannot generate active characters so we filter that: at some future stage that may change: the slightly odd ordering of auxiliaries reflects that.

```

10577 \group_begin:
10578   \char_set_catcode_active:N ^^L
10579   \cs_set:Npn ^^L { }
10580   \char_set_catcode_other:n { 0 }
10581   \if_int_odd:w 0
10582     \sys_if_engine luatex:T { 1 }
10583     \sys_if_engine xetex:T { 1 } \exp_stop_f:
10584     \cs_new:Npn \__char_generate_aux:nnw #1#2#3 \exp_end:
10585     {
10586       #3
10587       \exp_after:wN \exp_end:
10588       \tex_Ucharcat:D #1 \exp_stop_f: #2 \exp_stop_f:
10589     }
10590   \cs_if_exist:NF \tex_expanded:D
10591   {
10592     \cs_new_eq:NN \__char_generate_auxii:nnw \__char_generate_aux:nnw
10593     \cs_gset:Npn \__char_generate_aux:nnw #1#2#3 \exp_end:
10594     {
10595       #3
10596       \if_int_compare:w #2 = 13 \exp_stop_f:
10597         \__kernel_msg_expandable_error:nn { kernel } { char-active }
10598       \else:
10599         \__char_generate_auxii:nnw {#1} {#2}
10600       \fi:
10601       \exp_end:
10602     }
10603   }
10604 \else:

```

For engines where \Ucharcat isn't available or emulated, we have to work in macros, and cover only the 8-bit range. The first stage is to build up a `tl` containing `^^@` with each category code that can be accessed in this way, with an error set up for the other cases. This is all done such that it can be quickly accessed using a `\if_case:w` low-level

conditional. There are a few things to notice here. As `^^L` is `\outer` we need to locally set it to avoid a problem. To get open/close braces into the list, they are set up using `\if_false:` pairing and are then x-type expanded together into the desired form.

```

10605 \tl_set:Nn \l__char_tmp_tl { \exp_not:N \or: }
10606 \char_set_catcode_group_begin:n { 0 } % {
10607 \tl_put_right:Nn \l__char_tmp_tl { ^^@ \if_false: } }
10608 \char_set_catcode_group_end:n { 0 }
10609 \tl_put_right:Nn \l__char_tmp_tl { { \fi: \exp_not:N \or: ^^@ } % }
10610 \__kernel_tl_set:Nx \l__char_tmp_tl { \l__char_tmp_tl }
10611 \char_set_catcode_math_toggle:n { 0 }
10612 \tl_put_right:Nn \l__char_tmp_tl { \or: ^^@ }
10613 \char_set_catcode_alignment:n { 0 }
10614 \tl_put_right:Nn \l__char_tmp_tl { \or: ^^@ }
10615 \tl_put_right:Nn \l__char_tmp_tl { \or: }
10616 \char_set_catcode_parameter:n { 0 }
10617 \tl_put_right:Nn \l__char_tmp_tl { \or: ^^@ }
10618 \char_set_catcode_math_superscript:n { 0 }
10619 \tl_put_right:Nn \l__char_tmp_tl { \or: ^^@ }
10620 \char_set_catcode_math_subscript:n { 0 }
10621 \tl_put_right:Nn \l__char_tmp_tl { \or: ^^@ }
10622 \tl_put_right:Nn \l__char_tmp_tl { \or: }

```

For making spaces, there needs to be an o-type expansion of a `\use:n` (or some other tokenization) to avoid dropping the space.

```

10623 \char_set_catcode_space:n { 0 }
10624 \tl_put_right:No \l__char_tmp_tl { \use:n { \or: } ^^@ }
10625 \char_set_catcode_letter:n { 0 }
10626 \tl_put_right:Nn \l__char_tmp_tl { \or: ^^@ }
10627 \char_set_catcode_other:n { 0 }
10628 \tl_put_right:Nn \l__char_tmp_tl { \or: ^^@ }
10629 \char_set_catcode_active:n { 0 }
10630 \tl_put_right:Nn \l__char_tmp_tl { \or: ^^@ }

```

Convert the above temporary list into a series of constant token lists, one for each character code, using `\tex_lowercase:D` to convert `^^@` in each case. The x-type expansion ensures that `\tex_lowercase:D` receives the contents of the token list. `^^L` is awkward hence this is done in three parts: up to `^^L`, `^^L` itself and above `^^L`. Notice that at this stage `^^@` is active.

```

10631 \cs_set_protected:Npn \__char_tmp:n #1
10632 {
10633   \char_set_lccode:nn { 0 } {#1}
10634   \char_set_lccode:nn { 32 } {#1}
10635   \exp_args:Nx \tex_lowercase:D
10636   {
10637     \tl_const:Nn
10638       \exp_not:c { c__char_ \__char_int_to_roman:w #1 _tl }
10639       { \exp_not:o \l__char_tmp_tl }
10640   }
10641 }
10642 \int_step_function:nnN { 0 } { 11 } \__char_tmp:n
10643 \group_begin:
10644   \tl_replace_once:Nnn \l__char_tmp_tl { ^^@ } { \ERROR }
10645   \__char_tmp:n { 12 }
10646 \group_end:

```

```
10647 \int_step_function:nnN { 13 } { 255 } \__char_tmp:n
```

As T_EX is very unhappy if it finds an alignment character inside a primitive `\halign` even when skipping false branches, some precautions are required. T_EX is happy if the token is hidden between braces within `\if_false: ... \fi:`.

```
10648 \cs_new:Npn \__char_generate_aux:nnw #1#2#3 \exp_end:
10649 {
10650   #3
10651   \if_false: { \fi:
10652     \exp_after:wN \exp_after:wN
10653     \exp_after:wN \exp_end:
10654     \exp_after:wN \exp_after:wN
10655     \if_case:w #2
10656       \exp_last_unbraced:Nv \exp_stop_f:
10657       { c__char_ \__char_int_to_roman:w #1 _tl }
10658     \or: }
10659     \fi:
10660   }
10661 \fi:
10662 \group_end:
```

(End definition for `\char_generate:nn` and others. This function is documented on page 130.)

`\char_to_utfviii_bytes:n`

This code converts a codepoint into the correct UTF-8 representation. In terms of the algorithm itself, see <https://en.wikipedia.org/wiki/UTF-8> for the octet pattern.

```
\__char_to_utfviii_bytes_auxi:n
\__char_to_utfviii_bytes_auxii:Nnn
\__char_to_utfviii_bytes_auxiii:n
\__char_to_utfviii_bytes_outputi:nw
\__char_to_utfviii_bytes_outputii:nw
\__char_to_utfviii_bytes_outputiii:nw
\__char_to_utfviii_bytes_outputiv:nw
\__char_to_utfviii_bytes_output:nnn
\__char_to_utfviii_bytes_output:fn
\__char_to_utfviii_bytes_end:
10663 \cs_new:Npn \char_to_utfviii_bytes:n #1
10664 {
10665   \exp_args:Nf \__char_to_utfviii_bytes_auxi:n
10666   { \int_eval:n {#1} }
10667 }
10668 \cs_new:Npn \__char_to_utfviii_bytes_auxi:n #1
10669 {
10670   \if_int_compare:w #1 > "80 \exp_stop_f:
10671   \if_int_compare:w #1 < "800 \exp_stop_f:
10672     \__char_to_utfviii_bytes_outputi:nw
10673     { \__char_to_utfviii_bytes_auxii:Nnn C {#1} { 64 } }
10674     \__char_to_utfviii_bytes_outputii:nw
10675     { \__char_to_utfviii_bytes_auxiii:n {#1} }
10676   \else:
10677     \if_int_compare:w #1 < "10000 \exp_stop_f:
10678     \__char_to_utfviii_bytes_outputi:nw
10679     { \__char_to_utfviii_bytes_auxii:Nnn E {#1} { 64 * 64 } }
10680     \__char_to_utfviii_bytes_outputii:nw
10681     {
10682       \__char_to_utfviii_bytes_auxiii:n
10683       { \int_div_truncate:nn {#1} { 64 } }
10684     }
10685     \__char_to_utfviii_bytes_outputiii:nw
10686     { \__char_to_utfviii_bytes_auxiii:n {#1} }
10687   \else:
10688     \__char_to_utfviii_bytes_outputi:nw
10689     {
10690       \__char_to_utfviii_bytes_auxii:Nnn F
10691       {#1} { 64 * 64 * 64 }

```

```

10692     }
10693     \__char_to_utfviii_bytes_outputii:nw
10694     {
10695         \__char_to_utfviii_bytes_auxiii:n
10696         { \int_div_truncate:nn {#1} { 64 * 64 } }
10697     }
10698     \__char_to_utfviii_bytes_outputiii:nw
10699     {
10700         \__char_to_utfviii_bytes_auxiii:n
10701         { \int_div_truncate:nn {#1} { 64 } }
10702     }
10703     \__char_to_utfviii_bytes_outputiv:nw
10704     { \__char_to_utfviii_bytes_auxiii:n {#1} }
10705     \fi:
10706     \fi:
10707     \else:
10708         \__char_to_utfviii_bytes_outputi:nw {#1}
10709     \fi:
10710     \__char_to_utfviii_bytes_end: { } { } { } { }
10711 }
10712 \cs_new:Npn \__char_to_utfviii_bytes_auxii:Nnn #1#2#3
10713 { "#10 + \int_div_truncate:nn {#2} {#3} }
10714 \cs_new:Npn \__char_to_utfviii_bytes_auxiii:n #1
10715 { \int_mod:nn {#1} { 64 } + 128 }
10716 \cs_new:Npn \__char_to_utfviii_bytes_outputi:nw
10717 #1 #2 \__char_to_utfviii_bytes_end: #3
10718 { \__char_to_utfviii_bytes_output:fnn { \int_eval:n {#1} } { } {#2} }
10719 \cs_new:Npn \__char_to_utfviii_bytes_outputii:nw
10720 #1 #2 \__char_to_utfviii_bytes_end: #3#4
10721 { \__char_to_utfviii_bytes_output:fnn { \int_eval:n {#1} } { {#3} } {#2} }
10722 \cs_new:Npn \__char_to_utfviii_bytes_outputiii:nw
10723 #1 #2 \__char_to_utfviii_bytes_end: #3#4#5
10724 {
10725     \__char_to_utfviii_bytes_output:fnn
10726     { \int_eval:n {#1} } { {#3} {#4} } {#2}
10727 }
10728 \cs_new:Npn \__char_to_utfviii_bytes_outputiv:nw
10729 #1 #2 \__char_to_utfviii_bytes_end: #3#4#5#6
10730 {
10731     \__char_to_utfviii_bytes_output:fnn
10732     { \int_eval:n {#1} } { {#3} {#4} {#5} } {#2}
10733 }
10734 \cs_new:Npn \__char_to_utfviii_bytes_output:nnn #1#2#3
10735 {
10736     #3
10737     \__char_to_utfviii_bytes_end: #2 {#1}
10738 }
10739 \cs_generate_variant:Nn \__char_to_utfviii_bytes_output:nnn { f }
10740 \cs_new:Npn \__char_to_utfviii_bytes_end: { }

```

(End definition for `\char_to_utfviii_bytes:n` and others. This function is documented on page 276.)

```

\char_to_nfd:N Look up any NFD and recursively produce the result.
\__char_to_nfd:n
10741 \cs_new:Npn \char_to_nfd:N #1
\__char_to_nfd:Nw

```

```

10742 {
10743   \cs_if_exist:cTF { c__char_nfd_ \token_to_str:N #1 _ t1 }
10744   {
10745     \exp_after:wN \exp_after:wN \exp_after:wN \__char_to_nfd:Nw
10746     \exp_after:wN \exp_after:wN \exp_after:wN #1
10747     \cs:w c__char_nfd_ \token_to_str:N #1 _ t1 \cs_end:
10748     \s__char_stop
10749   }
10750   { \exp_not:n {#1} }
10751 }
10752 \cs_set_eq:NN \__char_to_nfd:n \char_to_nfd:N
10753 \cs_new:Npn \__char_to_nfd:Nw #1#2#3 \s__char_stop
10754 {
10755   \exp_args:Ne \__char_to_nfd:n
10756   { \char_generate:nn { '#2 } { \__char_change_case_catcode:N #1 } }
10757   \tl_if_blank:nF {#3}
10758   {
10759     \exp_args:Ne \__char_to_nfd:n
10760     { \char_generate:nn { '#3 } { \char_value_catcode:n { '#3 } } }
10761   }
10762 }

```

(End definition for `\char_to_nfd:N`, `__char_to_nfd:n`, and `__char_to_nfd:Nw`. This function is documented on page 276.)

`\char_lowercase:N` To ensure that the category codes produced are predictable, every character is re-generated even if it is otherwise unchanged. This makes life a little interesting when we might have multiple output characters: we have to grab each of them and case change them in reverse order to maintain f-type expandability.

```

\__char_change_case:nNN 10763 \cs_new:Npn \char_lowercase:N #1
\__char_change_case:nN 10764 { \__char_change_case:nNN { lower } \char_value_lccode:n #1 }
\__char_change_case_multi:nN 10765 \cs_new:Npn \char_uppercase:N #1
\__char_change_case_multi:vN 10766 { \__char_change_case:nNN { upper } \char_value_uccode:n #1 }
\__char_change_case_multi:NNNw 10767 \cs_new:Npn \char_titlecase:N #1
\__char_change_case:NNN 10768 {
\__char_change_case:NNNN 10769   \tl_if_exist:cTF { c__char_titlecase_ \token_to_str:N #1 _t1 }
\__char_change_case:NN 10770   {
\__char_change_case_catcode:N 10771     \__char_change_case_multi:vN
10772     { c__char_titlecase_ \token_to_str:N #1 _t1 } #1
10773   }
10774   { \char_uppercase:N #1 }
10775 }
\char_str_lowercase:N 10776 \cs_new:Npn \char_foldcase:N #1
\char_str_uppercase:N 10777 { \__char_change_case:nNN { fold } \char_value_lccode:n #1 }
\char_str_titlecase:N 10778 \cs_new:Npn \__char_change_case:nNN #1#2#3
\char_str_foldcase:N 10779 {
\__char_str_change_case:nNN 10780   \tl_if_exist:cTF { c__char_ #1 case_ \token_to_str:N #3 _t1 }
\__char_str_change_case:nN 10781   {
10782     \__char_change_case_multi:vN
10783     { c__char_ #1 case_ \token_to_str:N #3 _t1 } #3
10784   }
10785   { \exp_args:Nf \__char_change_case:nN { #2 { '#3 } } #3 }
10786 }
10787 \cs_new:Npn \__char_change_case:nN #1#2

```

```

10788 {
10789   \int_compare:nNnTF {#1} = 0
10790     { #2 }
10791     { \char_generate:nn {#1} { \__char_change_case_catcode:N #2 } }
10792 }
10793 \cs_new:Npn \__char_change_case_multi:nN #1#2
10794 { \__char_change_case_multi:NNNNw #2 #1 \q__char_no_value \q__char_no_value \s__char_stop
10795 \cs_generate_variant:Nn \__char_change_case_multi:nN { v }
10796 \cs_new:Npn \__char_change_case_multi:NNNNw #1#2#3#4#5 \s__char_stop
10797 {
10798   \__char_quark_if_no_value:NTF #4
10799   {
10800     \__char_quark_if_no_value:NTF #3
10801     { \__char_change_case:NN #1 #2 }
10802     { \__char_change_case:NNN #1 #2#3 }
10803   }
10804   { \__char_change_case:NNNN #1 #2#3#4 }
10805 }
10806 \cs_new:Npn \__char_change_case:NNN #1#2#3
10807 {
10808   \exp_args:Nnf \use:nn
10809   { \__char_change_case:NN #1 #2 }
10810   { \__char_change_case:NN #1 #3 }
10811 }
10812 \cs_new:Npn \__char_change_case:NNNN #1#2#3#4
10813 {
10814   \exp_args:Nfff \use:nnn
10815   { \__char_change_case:NN #1 #2 }
10816   { \__char_change_case:NN #1 #3 }
10817   { \__char_change_case:NN #1 #4 }
10818 }
10819 \cs_new:Npn \__char_change_case:NN #1#2
10820 { \char_generate:nn { '#2 } { \__char_change_case_catcode:N #1 } }
10821 \cs_new:Npn \__char_change_case_catcode:N #1
10822 {
10823   \if_catcode:w \exp_not:N #1 \c_math_toggle_token
10824   3
10825   \else:
10826     \if_catcode:w \exp_not:N #1 \c_alignment_token
10827     4
10828     \else:
10829       \if_catcode:w \exp_not:N #1 \c_math_superscript_token
10830       7
10831       \else:
10832         \if_catcode:w \exp_not:N #1 \c_math_subscript_token
10833         8
10834         \else:
10835           \if_catcode:w \exp_not:N #1 \c_space_token
10836           10
10837           \else:
10838             \if_catcode:w \exp_not:N #1 \c_catcode_letter_token
10839             11
10840             \else:
10841               \if_catcode:w \exp_not:N #1 \c_catcode_other_token

```

```

10842             12
10843             \else:
10844             13
10845             \fi:
10846             \fi:
10847             \fi:
10848             \fi:
10849             \fi:
10850             \fi:
10851             \fi:
10852         }

```

Same story for the string version, except category code is easier to follow. This of course makes this version significantly faster.

```

10853 \cs_new:Npn \char_str_lowercase:N #1
10854 { \__char_str_change_case:nNN { lower } \char_value_lccode:n #1 }
10855 \cs_new:Npn \char_str_uppercase:N #1
10856 { \__char_str_change_case:nNN { upper } \char_value_uccode:n #1 }
10857 \cs_new:Npn \char_str_titlecase:N #1
10858 {
10859     \tl_if_exist:cTF { c__char_titlecase_ \token_to_str:N #1 _tl }
10860     { \tl_to_str:c { c__char_titlecase_ \token_to_str:N #1 _tl } }
10861     { \char_str_uppercase:N #1 }
10862 }
10863 \cs_new:Npn \char_str_foldcase:N #1
10864 { \__char_str_change_case:nNN { fold } \char_value_lccode:n #1 }
10865 \cs_new:Npn \__char_str_change_case:nNN #1#2#3
10866 {
10867     \tl_if_exist:cTF { c__char_ #1 case_ \token_to_str:N #3 _tl }
10868     { \tl_to_str:c { c__char_ #1 case_ \token_to_str:N #3 _tl } }
10869     { \exp_args:Nf \__char_str_change_case:nN { #2 { '#3 } } #3 }
10870 }
10871 \cs_new:Npn \__char_str_change_case:nN #1#2
10872 {
10873     \int_compare:nNnTF {#1} = 0
10874     { \tl_to_str:n {#2} }
10875     { \char_generate:nn {#1} { 12 } }
10876 }
10877 \bool_lazy_or:nnF
10878 { \cs_if_exist_p:N \tex_luatexversion:D }
10879 { \cs_if_exist_p:N \tex_XeTeXversion:D }
10880 {
10881     \cs_set:Npn \__char_str_change_case:nN #1#2
10882     { \tl_to_str:n {#2} }
10883 }

```

(End definition for `\char_lowercase:N` and others. These functions are documented on page 130.)

`\c_catcode_other_space_tl` Create a space with category code 12: an “other” space.

```

10884 \tl_const:Nx \c_catcode_other_space_tl { \char_generate:nn { '\ } { 12 } }

```

(End definition for `\c_catcode_other_space_tl`. This function is documented on page 130.)

17.4 Generic tokens

10885 <@@=token>

\s__token_mark Internal scan marks.

10886 \scan_new:N \s__token_mark

10887 \scan_new:N \s__token_stop

(End definition for \s__token_mark and \s__token_stop.)

\token_to_meaning:N These are all defined in l3basics, as they are needed “early”. This is just a reminder!

\token_to_meaning:c

\token_to_str:N

\token_to_str:c

(End definition for \token_to_meaning:N and \token_to_str:N. These functions are documented on page 134.)

\c_group_begin_token

\c_group_end_token

\c_math_toggle_token

\c_alignment_token

\c_parameter_token

We define these useful tokens. For the brace and space tokens things have to be done by hand: the formal argument spec. for \cs_new_eq:NN does not cover them so we do things by hand. (As currently coded it would *work* with \cs_new_eq:NN but that’s not really a great idea to show off: we want people to stick to the defined interfaces and that includes us.) So that these few odd names go into the log when appropriate there is a need to hand-apply the __kernel_chk_if_free_cs:N check.

\c_math_superscript_token

\c_math_subscript_token

\c_space_token

\c_catcode_letter_token

\c_catcode_other_token

10888 \group_begin:

10889 __kernel_chk_if_free_cs:N \c_group_begin_token

10890 \tex_global:D \tex_let:D \c_group_begin_token {

10891 __kernel_chk_if_free_cs:N \c_group_end_token

10892 \tex_global:D \tex_let:D \c_group_end_token }

10893 \char_set_catcode_math_toggle:N *

10894 \cs_new_eq:NN \c_math_toggle_token *

10895 \char_set_catcode_alignment:N *

10896 \cs_new_eq:NN \c_alignment_token *

10897 \cs_new_eq:NN \c_parameter_token #

10898 \cs_new_eq:NN \c_math_superscript_token ^

10899 \char_set_catcode_math_subscript:N *

10900 \cs_new_eq:NN \c_math_subscript_token *

10901 __kernel_chk_if_free_cs:N \c_space_token

10902 \use:n { \tex_global:D \tex_let:D \c_space_token = ~ } ~

10903 \cs_new_eq:NN \c_catcode_letter_token a

10904 \cs_new_eq:NN \c_catcode_other_token 1

10905 \group_end:

(End definition for \c_group_begin_token and others. These functions are documented on page 134.)

\c_catcode_active_tl Not an implicit token!

10906 \group_begin:

10907 \char_set_catcode_active:N *

10908 \tl_const:Nn \c_catcode_active_tl { \exp_not:N * }

10909 \group_end:

(End definition for \c_catcode_active_tl. This variable is documented on page 134.)

17.5 Token conditionals

`\token_if_group_begin_p:N` Check if token is a begin group token. We use the constant `\c_group_begin_token` for this.
`\token_if_group_begin:N \underline{TF}`

```
10910 \prg_new_conditional:Npnn \token_if_group_begin:N #1 { p , T , F , TF }
10911 {
10912     \if_catcode:w \exp_not:N #1 \c_group_begin_token
10913     \prg_return_true: \else: \prg_return_false: \fi:
10914 }
```

(End definition for `\token_if_group_begin:N \underline{TF}` . This function is documented on page 135.)

`\token_if_group_end_p:N` Check if token is a end group token. We use the constant `\c_group_end_token` for this.
`\token_if_group_end:N \underline{TF}`

```
10915 \prg_new_conditional:Npnn \token_if_group_end:N #1 { p , T , F , TF }
10916 {
10917     \if_catcode:w \exp_not:N #1 \c_group_end_token
10918     \prg_return_true: \else: \prg_return_false: \fi:
10919 }
```

(End definition for `\token_if_group_end:N \underline{TF}` . This function is documented on page 135.)

`\token_if_math_toggle_p:N` Check if token is a math shift token. We use the constant `\c_math_toggle_token` for this.
`\token_if_math_toggle:N \underline{TF}`

```
10920 \prg_new_conditional:Npnn \token_if_math_toggle:N #1 { p , T , F , TF }
10921 {
10922     \if_catcode:w \exp_not:N #1 \c_math_toggle_token
10923     \prg_return_true: \else: \prg_return_false: \fi:
10924 }
```

(End definition for `\token_if_math_toggle:N \underline{TF}` . This function is documented on page 135.)

`\token_if_alignment_p:N` Check if token is an alignment tab token. We use the constant `\c_alignment_token` for this.
`\token_if_alignment:N \underline{TF}`

```
10925 \prg_new_conditional:Npnn \token_if_alignment:N #1 { p , T , F , TF }
10926 {
10927     \if_catcode:w \exp_not:N #1 \c_alignment_token
10928     \prg_return_true: \else: \prg_return_false: \fi:
10929 }
```

(End definition for `\token_if_alignment:N \underline{TF}` . This function is documented on page 135.)

`\token_if_parameter_p:N` Check if token is a parameter token. We use the constant `\c_parameter_token` for this.
`\token_if_parameter:N \underline{TF}` We have to trick T_EX a bit to avoid an error message: within a group we prevent `\c_parameter_token` from behaving like a macro parameter character. The definitions of `\prg_new_conditional:Npnn` are global, so they remain after the group.

```
10930 \group_begin:
10931 \cs_set_eq:NN \c_parameter_token \scan_stop:
10932 \prg_new_conditional:Npnn \token_if_parameter:N #1 { p , T , F , TF }
10933 {
10934     \if_catcode:w \exp_not:N #1 \c_parameter_token
10935     \prg_return_true: \else: \prg_return_false: \fi:
10936 }
10937 \group_end:
```

(End definition for `\token_if_parameter:N \underline{TF}` . This function is documented on page 135.)

`\token_if_math_superscript_p:N` Check if token is a math superscript token. We use the constant `\c_math_superscript_token` for this.

`\token_if_math_superscript:N \underline{TF}`

```

10938 \prg_new_conditional:Npnn \token_if_math_superscript:N #1
10939 { p , T , F , TF }
10940 {
10941     \if_catcode:w \exp_not:N #1 \c_math_superscript_token
10942     \prg_return_true: \else: \prg_return_false: \fi:
10943 }

```

(End definition for `\token_if_math_superscript:N \underline{TF}` . This function is documented on page 135.)

`\token_if_math_subscript_p:N` Check if token is a math subscript token. We use the constant `\c_math_subscript_token` for this.

`\token_if_math_subscript:N \underline{TF}`

```

10944 \prg_new_conditional:Npnn \token_if_math_subscript:N #1 { p , T , F , TF }
10945 {
10946     \if_catcode:w \exp_not:N #1 \c_math_subscript_token
10947     \prg_return_true: \else: \prg_return_false: \fi:
10948 }

```

(End definition for `\token_if_math_subscript:N \underline{TF}` . This function is documented on page 135.)

`\token_if_space_p:N` Check if token is a space token. We use the constant `\c_space_token` for this.

`\token_if_space:N \underline{TF}`

```

10949 \prg_new_conditional:Npnn \token_if_space:N #1 { p , T , F , TF }
10950 {
10951     \if_catcode:w \exp_not:N #1 \c_space_token
10952     \prg_return_true: \else: \prg_return_false: \fi:
10953 }

```

(End definition for `\token_if_space:N \underline{TF}` . This function is documented on page 135.)

`\token_if_letter_p:N` Check if token is a letter token. We use the constant `\c_catcode_letter_token` for this.

`\token_if_letter:N \underline{TF}`

```

10954 \prg_new_conditional:Npnn \token_if_letter:N #1 { p , T , F , TF }
10955 {
10956     \if_catcode:w \exp_not:N #1 \c_catcode_letter_token
10957     \prg_return_true: \else: \prg_return_false: \fi:
10958 }

```

(End definition for `\token_if_letter:N \underline{TF}` . This function is documented on page 136.)

`\token_if_other_p:N` Check if token is an other char token. We use the constant `\c_catcode_other_token` for this.

`\token_if_other:N \underline{TF}`

```

10959 \prg_new_conditional:Npnn \token_if_other:N #1 { p , T , F , TF }
10960 {
10961     \if_catcode:w \exp_not:N #1 \c_catcode_other_token
10962     \prg_return_true: \else: \prg_return_false: \fi:
10963 }

```

(End definition for `\token_if_other:N \underline{TF}` . This function is documented on page 136.)

`\token_if_active_p:N` Check if token is an active char token. We use the constant `\c_catcode_active_tl` for this. A technical point is that `\c_catcode_active_tl` is in fact a macro expanding to `\exp_not:N *`, where `*` is active.

`\token_if_active:N \underline{TF}`

```

10964 \prg_new_conditional:Npnn \token_if_active:N #1 { p , T , F , TF }
10965 {
10966     \if_catcode:w \exp_not:N #1 \c_catcode_active_tl
10967     \prg_return_true: \else: \prg_return_false: \fi:
10968 }

```

(End definition for `\token_if_active:NTF`. This function is documented on page 136.)

`\token_if_eq_meaning_p:NN` Check if the tokens #1 and #2 have same meaning.

`\token_if_eq_meaning:NNTF`

```

10969 \prg_new_conditional:Npnn \token_if_eq_meaning:NN #1#2 { p , T , F , TF }
10970 {
10971   \if_meaning:w #1 #2
10972   \prg_return_true: \else: \prg_return_false: \fi:
10973 }

```

(End definition for `\token_if_eq_meaning:NNTF`. This function is documented on page 136.)

`\token_if_eq_catcode_p:NN` Check if the tokens #1 and #2 have same category code.

`\token_if_eq_catcode:NNTF`

```

10974 \prg_new_conditional:Npnn \token_if_eq_catcode:NN #1#2 { p , T , F , TF }
10975 {
10976   \if_catcode:w \exp_not:N #1 \exp_not:N #2
10977   \prg_return_true: \else: \prg_return_false: \fi:
10978 }

```

(End definition for `\token_if_eq_catcode:NNTF`. This function is documented on page 136.)

`\token_if_eq_charcode_p:NN` Check if the tokens #1 and #2 have same character code.

`\token_if_eq_charcode:NNTF`

```

10979 \prg_new_conditional:Npnn \token_if_eq_charcode:NN #1#2 { p , T , F , TF }
10980 {
10981   \if_charcode:w \exp_not:N #1 \exp_not:N #2
10982   \prg_return_true: \else: \prg_return_false: \fi:
10983 }

```

(End definition for `\token_if_eq_charcode:NNTF`. This function is documented on page 136.)

`\token_if_macro_p:N` When a token is a macro, `\token_to_meaning:N` always outputs something like
`\token_if_macro:NNTF` `\long macro:#1->#1` so we could naively check to see if the meaning contains `->`.
`__token_if_macro_p:w` However, this can fail the five `\...mark` primitives, whose meaning has the form `\...mark:<user material>`. The problem is that the `<user material>` can contain `->`.

However, only characters, macros, and marks can contain the colon character. The idea is thus to grab until the first `:`, and analyse what is left. However, macros can have any combination of `\long`, `\protected` or `\outer` (not used in L^AT_EX3) before the string `macro:.` We thus only select the part of the meaning between the first `ma` and the first following `:`. If this string is `cro`, then we have a macro. If the string is `rk`, then we have a mark. The string can also be `cro parameter character` for a colon with a weird category code (namely the usual category code of `#`). Otherwise, it is empty.

This relies on the fact that `\long`, `\protected`, `\outer` cannot contain `ma`, regardless of the escape character, even if the escape character is `m...`

Both `ma` and `:` must be of category code 12 (other), so are detokenized.

```

10984 \use:x
10985 {
10986   \prg_new_conditional:Npnn \exp_not:N \token_if_macro:N ##1
10987   { p , T , F , TF }
10988   {
10989     \exp_not:N \exp_after:wN \exp_not:N \__token_if_macro_p:w
10990     \exp_not:N \token_to_meaning:N ##1 \tl_to_str:n { ma : }
10991     \s__token_stop
10992   }
10993   \cs_new:Npn \exp_not:N \__token_if_macro_p:w

```

```

10994     ##1 \tl_to_str:n { ma } ##2 \c_colon_str ##3 \s__token_stop
10995   }
10996   {
10997     \str_if_eq:nnTF { #2 } { cro }
10998     { \prg_return_true: }
10999     { \prg_return_false: }
11000   }

```

(End definition for `\token_if_macro:NTF` and `__token_if_macro_p:w`. This function is documented on page 136.)

`\token_if_cs_p:N` Check if token has same catcode as a control sequence. This follows the same pattern as `\token_if_letter:N` etc. We use `\scan_stop:` for this.

`\token_if_cs:NTF`

```

11001 \prg_new_conditional:Npnn \token_if_cs:N #1 { p , T , F , TF }
11002 {
11003   \if_catcode:w \exp_not:N #1 \scan_stop:
11004   \prg_return_true: \else: \prg_return_false: \fi:
11005 }

```

(End definition for `\token_if_cs:NTF`. This function is documented on page 136.)

`\token_if_expandable_p:N` Check if token is expandable. We use the fact that T_EX temporarily converts `\exp_not:N` $\langle token \rangle$ into `\scan_stop:` if $\langle token \rangle$ is expandable. An undefined token is not considered as expandable. No problem nesting the conditionals, since the third #1 is only skipped if it is non-expandable (hence not part of T_EX’s conditional apparatus).

`\token_if_expandable:NTF`

```

11006 \prg_new_conditional:Npnn \token_if_expandable:N #1 { p , T , F , TF }
11007 {
11008   \exp_after:wN \if_meaning:w \exp_not:N #1 #1
11009   \prg_return_false:
11010   \else:
11011     \if_cs_exist:N #1
11012     \prg_return_true:
11013     \else:
11014       \prg_return_false:
11015     \fi:
11016   \fi:
11017 }

```

(End definition for `\token_if_expandable:NTF`. This function is documented on page 136.)

`__token_delimit_by_char:w` These auxiliary functions are used below to define some conditionals which detect whether the `\meaning` of their argument begins with a particular string. Each auxiliary takes an argument delimited by a string, a second one delimited by `\s__token_stop`, and returns the first one and its delimiter. This result is eventually compared to another string. Note that the “font” auxiliary is delimited by a space followed by “font”. This avoids an unnecessary check for the `\font` primitive below.

`__token_delimit_by_count:w`
`__token_delimit_by_dimen:w`
`__token_delimit_by_font:w`
`__token_delimit_by_macro:w`
`__token_delimit_by_muskip:w`
`__token_delimit_by_skip:w`
`__token_delimit_by_toks:w`

```

11018 \group_begin:
11019 \cs_set_protected:Npn \__token_tmp:w #1
11020 {
11021   \use:x
11022   {
11023     \cs_new:Npn \exp_not:c { __token_delimit_by_ #1 :w }
11024     #####1 \tl_to_str:n {#1} #####2 \s__token_stop
11025     { #####1 \tl_to_str:n {#1} }

```

```

11026     }
11027   }
11028   \__token_tmp:w { char" }
11029   \__token_tmp:w { count }
11030   \__token_tmp:w { dimen }
11031   \__token_tmp:w { ~ font }
11032   \__token_tmp:w { macro }
11033   \__token_tmp:w { muskip }
11034   \__token_tmp:w { skip }
11035   \__token_tmp:w { toks }
11036 \group_end:

```

(End definition for `__token_delimit_by_char:w` and others.)

Each of these conditionals tests whether its argument's `\meaning` starts with a given string. This is essentially done by having an auxiliary grab an argument delimited by the string and testing whether the argument was empty. Of course, a copy of this string must first be added to the end of the `\meaning` to avoid a runaway argument in case it does not contain the string. Two complications arise. First, the escape character is not fixed, and cannot be included in the delimiter of the auxiliary function (this function cannot be defined on the fly because tests must remain expandable): instead the first argument of the auxiliary (plus the delimiter to avoid complications with trailing spaces) is compared using `\str_if_eq:eeTF` to the result of applying `\token_to_str:N` to a control sequence. Second, the `\meaning` of primitives such as `\dimen` or `\dimendef` starts in the same way as registers such as `\dimen123`, so they must be tested for.

Characters used as delimiters must have catcode 12 and are obtained through `\tl_to_str:n`. This requires doing all definitions within `x`-expansion. The temporary function `__token_tmp:w` used to define each conditional receives three arguments: the name of the conditional, the auxiliary's delimiter (also used to name the auxiliary), and the string to which one compares the auxiliary's result. Note that the `\meaning` of a protected long macro starts with `\protected\long macro`, with no space after `\protected` but a space after `\long`, hence the mixture of `\token_to_str:N` and `\tl_to_str:n`.

For the first six conditionals, `\cs_if_exist:cT` turns out to be `false` (thanks to the leading space for `font`), and the code boils down to a string comparison between the result of the auxiliary on the `\meaning` of the conditional's argument `####1`, and `#3`. Both are evaluated at run-time, as this is important to get the correct escape character.

The other five conditionals have additional code that compares the argument `####1` to two `TeX` primitives which would wrongly be recognized as registers otherwise. Despite using `TeX`'s primitive conditional construction, this does not break when `####1` is itself a conditional, because branches of the conditionals are only skipped if `####1` is one of the two primitives that are tested for (which are not `TeX` conditionals).

```

11037 \group_begin:
11038 \cs_set_protected:Npn \__token_tmp:w #1#2#3
11039 {
11040   \use:x
11041   {
11042     \prg_new_conditional:Npnn \exp_not:c { token_if_ #1 :N } ####1
11043     { p , T , F , TF }
11044     {
11045       \cs_if_exist:cT { tex_ #2 :D }
11046       {
11047         \exp_not:N \if_meaning:w ####1 \exp_not:c { tex_ #2 :D }

```

```

11048         \exp_not:N \prg_return_false:
11049         \exp_not:N \else:
11050         \exp_not:N \if_meaning:w #####1 \exp_not:c { tex_ #2 def:D }
11051         \exp_not:N \prg_return_false:
11052         \exp_not:N \else:
11053     }
11054     \exp_not:N \str_if_eq:eeTF
11055     {
11056         \exp_not:N \exp_after:wN
11057         \exp_not:c { __token_delimit_by_ #2 :w }
11058         \exp_not:N \token_to_meaning:N #####1
11059         ? \tl_to_str:n {#2} \s__token_stop
11060     }
11061     { \exp_not:n {#3} }
11062     { \exp_not:N \prg_return_true: }
11063     { \exp_not:N \prg_return_false: }
11064     \cs_if_exist:cT { tex_ #2 :D }
11065     {
11066         \exp_not:N \fi:
11067         \exp_not:N \fi:
11068     }
11069 }
11070 }
11071 }
11072 \__token_tmp:w { chardef } { char" } { \token_to_str:N \char" }
11073 \__token_tmp:w { mathchardef } { char" } { \token_to_str:N \mathchar" }
11074 \__token_tmp:w { long_macro } { macro } { \tl_to_str:n { \long } macro }
11075 \__token_tmp:w { protected_macro } { macro }
11076     { \tl_to_str:n { \protected } macro }
11077 \__token_tmp:w { protected_long_macro } { macro }
11078     { \token_to_str:N \protected \tl_to_str:n { \long } macro }
11079 \__token_tmp:w { font_selection } { ~ font } { select ~ font }
11080 \__token_tmp:w { dim_register } { dimen } { \token_to_str:N \dimen }
11081 \__token_tmp:w { int_register } { count } { \token_to_str:N \count }
11082 \__token_tmp:w { muskip_register } { muskip } { \token_to_str:N \muskip }
11083 \__token_tmp:w { skip_register } { skip } { \token_to_str:N \skip }
11084 \__token_tmp:w { toks_register } { toks } { \token_to_str:N \toks }
11085 \group_end:

```

(End definition for `\token_if_chardef:N` and others. These functions are documented on page 137.)

`\token_if_primitive_p:N`

`\token_if_primitive:N`

We filter out macros first, because they cause endless trouble later otherwise.

Primitives are almost distinguished by the fact that the result of `\token_to_meaning:N` is formed from letters only. Every other token has either a space (e.g., the letter A), a digit (e.g., `\count123`) or a double quote (e.g., `\char"A`).

Ten exceptions: on the one hand, `\tex_undefined:D` is not a primitive, but its meaning is undefined, only letters; on the other hand, `\space`, `\italiccorr`, `\hyphen`, `\firstmark`, `\topmark`, `\botmark`, `\splitfirstmark`, `\splitbotmark`, and `\nullfont` are primitives, but have non-letters in their meaning.

We start by removing the two first (non-space) characters from the meaning. This removes the escape character (which may be nonexistent depending on `\endlinechar`), and takes care of three of the exceptions: `\space`, `\italiccorr` and `\hyphen`, whose meaning is at most two characters. This leaves a string terminated by some `:`, and `\s__token_stop`.

The meaning of each one of the five `\...mark` primitives has the form $\langle letters \rangle : \langle user material \rangle$. In other words, the first non-letter is a colon. We remove everything after the first colon.

We are now left with a string, which we must analyze. For primitives, it contains only letters. For non-primitives, it contains either " , or a space, or a digit. Two exceptions remain: `\tex_undefined:D`, which is not a primitive, and `\nullfont`, which is a primitive.

Spaces cannot be grabbed in an undelimited way, so we check them separately. If there is a space, we test for `\nullfont`. Otherwise, we go through characters one by one, and stop at the first character less than 'A' (this is not quite a test for "only letters", but is close enough to work in this context). If this first character is : then we have a primitive, or `\tex_undefined:D`, and if it is " or a digit, then the token is not a primitive.

For LuaTeX we use a different implementation which just looks at the command code for the token and compaes it to a list of non-primitives. Again, `\nullfont` is a special case because it is the only primitive with the normally non-primitive `set_font` command code.

```

11086 \sys_if_engine luatex:TF
11087 {
11088   </tex>
11089   <*lua>
11090 do
11091   local get_next = token.get_next
11092   local get_command = token.get_command
11093   local get_index = token.get_index
11094   local get_mode = token.get_mode or token.get_index
11095   local cmd = token.command_id
11096   local set_font = cmd'get_font'
11097   local biggest_char = token.biggest_char()
11098
11099   local mode_below_biggest_char = {}
11100   local index_not_nil = {}
11101   local mode_not_null = {}
11102   local non_primitive = {
11103     [cmd'left_brace'] = true,
11104     [cmd'right_brace'] = true,
11105     [cmd'math_shift'] = true,
11106     [cmd'mac_param'] = mode_below_biggest_char,
11107     [cmd'sup_mark'] = true,
11108     [cmd'sub_mark'] = true,
11109     [cmd'endv'] = true,
11110     [cmd'spacer'] = true,
11111     [cmd'letter'] = true,
11112     [cmd'other_char'] = true,
11113     [cmd'tab_mark'] = mode_below_biggest_char,
11114     [cmd'char_given'] = true,
11115     [cmd'math_given'] = true,
11116     [cmd'xmath_given'] = true,
11117     [cmd'set_font'] = mode_not_null,
11118     [cmd'undefined_cs'] = true,
11119     [cmd'call'] = true,
11120     [cmd'long_call'] = true,
11121     [cmd'outer_call'] = true,

```

```

11122 [cmd'long_outer_call'] = true,
11123 [cmd'assign_glue'] = index_not_nil,
11124 [cmd'assign_mu_glue'] = index_not_nil,
11125 [cmd'assign_toks'] = index_not_nil,
11126 [cmd'assign_int'] = index_not_nil,
11127 [cmd'assign_attr'] = true,
11128 [cmd'assign_dimen'] = index_not_nil,
11129 }
11130
11131 luacmd("__token_if_primitive_lua:N", function()
11132   local tok = get_next()
11133   local is_non_primitive = non_primitive[get_command(tok)]
11134   return put_next(
11135     is_non_primitive == true
11136     and false_tok
11137   or is_non_primitive == nil
11138     and true_tok
11139   or is_non_primitive == mode_not_null
11140     and (get_mode(tok) == 0 and true_tok or false_tok)
11141   or is_non_primitive == index_not_nil
11142     and (get_index(tok) and false_tok or true_tok)
11143   or is_non_primitive == mode_below_biggest_char
11144     and (get_mode(tok) > biggest_char and true_tok or false_tok))
11145   end, "global")
11146 end
11147 </lua>
11148 <*tex>
11149 \prg_new_conditional:Npnn \token_if_primitive:N #1 { p , T , F , TF }
11150 {
11151   \__token_if_primitive_lua:N #1
11152 }
11153 }
11154 {
11155   \tex_chardef:D \c__token_A_int = 'A ~ %
11156   \use:x
11157   {
11158     \prg_new_conditional:Npnn \exp_not:N \token_if_primitive:N ##1
11159     { p , T , F , TF }
11160     {
11161       \exp_not:N \token_if_macro:NTF ##1
11162       \exp_not:N \prg_return_false:
11163       {
11164         \exp_not:N \exp_after:wN \exp_not:N \__token_if_primitive:NNw
11165         \exp_not:N \token_to_meaning:N ##1
11166         \tl_to_str:n { : : : } \s__token_stop ##1
11167       }
11168     }
11169     \cs_new:Npn \exp_not:N \__token_if_primitive:NNw
11170     ##1##2 ##3 \c_colon_str ##4 \s__token_stop
11171     {
11172       \exp_not:N \tl_if_empty:oTF
11173       { \exp_not:N \__token_if_primitive_space:w ##3 ~ }
11174       {
11175         \exp_not:N \__token_if_primitive_loop:N ##3

```



```

11176         \c_colon_str \s__token_stop
11177     }
11178     { \exp_not:N \__token_if_primitive_nullfont:N }
11179 }
11180 }
11181 \cs_new:Npn \__token_if_primitive_space:w #1 ~ { }
11182 \cs_new:Npn \__token_if_primitive_nullfont:N #1
11183 {
11184     \if_meaning:w \tex_nullfont:D #1
11185     \prg_return_true:
11186 }else:
11187     \prg_return_false:
11188 \fi:
11189 }
11190 \cs_new:Npn \__token_if_primitive_loop:N #1
11191 {
11192     \if_int_compare:w '#1 < \c__token_A_int %
11193     \exp_after:wN \__token_if_primitive:Nw
11194     \exp_after:wN #1
11195 }else:
11196     \exp_after:wN \__token_if_primitive_loop:N
11197 \fi:
11198 }
11199 \cs_new:Npn \__token_if_primitive:Nw #1 #2 \s__token_stop
11200 {
11201     \if:w : #1
11202     \exp_after:wN \__token_if_primitive_undefined:N
11203 }else:
11204     \prg_return_false:
11205     \exp_after:wN \use_none:n
11206 \fi:
11207 }
11208 \cs_new:Npn \__token_if_primitive_undefined:N #1
11209 {
11210     \if_cs_exist:N #1
11211     \prg_return_true:
11212 }else:
11213     \prg_return_false:
11214 \fi:
11215 }
11216 }

```

(End definition for `\token_if_primitive:NTF` and others. This function is documented on page 138.)

`\token_case_catcode:Nn` The aim here is to allow the case statement to be evaluated using a known number of expansion steps (two), and without needing to use an explicit “end of recursion” marker. That is achieved by using the test input as the final case, as this is always true. The trick is then to tidy up the output such that the appropriate case code plus either the true or false branch code is inserted.

`\token_case_catcode:NnTF`

`\token_case_charcode:Nn`

`\token_case_charcode:NnTF`

`\token_case_meaning:Nn`

`\token_case_meaning:NnTF`

```

11217 \cs_new:Npn \token_case_catcode:Nn #1#2
11218 { \exp:w \__token_case:NNnTF \token_if_eq_catcode:NNTF #1 {#2} { } { } }
11219 \cs_new:Npn \token_case_catcode:NnT #1#2#3
11220 { \exp:w \__token_case:NNnTF \token_if_eq_catcode:NNTF #1 {#2} {#3} { } }
11221 \cs_new:Npn \token_case_catcode:NnF #1#2

```

```

11222 { \exp:w \_token_case:NNnTF \token_if_eq_catcode:NNTF #1 {#2} { } }
11223 \cs_new:Npn \token_case_catcode:NnTF
11224 { \exp:w \_token_case:NNnTF \token_if_eq_catcode:NNTF }
11225 \cs_new:Npn \token_case_charcode:Nn #1#2
11226 { \exp:w \_token_case:NNnTF \token_if_eq_charcode:NNTF #1 {#2} { } { } }
11227 \cs_new:Npn \token_case_charcode:NnT #1#2#3
11228 { \exp:w \_token_case:NNnTF \token_if_eq_charcode:NNTF #1 {#2} {#3} { } }
11229 \cs_new:Npn \token_case_charcode:NnF #1#2
11230 { \exp:w \_token_case:NNnTF \token_if_eq_charcode:NNTF #1 {#2} { } }
11231 \cs_new:Npn \token_case_charcode:NnTF
11232 { \exp:w \_token_case:NNnTF \token_if_eq_charcode:NNTF }
11233 \cs_new:Npn \token_case_meaning:Nn #1#2
11234 { \exp:w \_token_case:NNnTF \token_if_eq_meaning:NNTF #1 {#2} { } { } }
11235 \cs_new:Npn \token_case_meaning:NnT #1#2#3
11236 { \exp:w \_token_case:NNnTF \token_if_eq_meaning:NNTF #1 {#2} {#3} { } }
11237 \cs_new:Npn \token_case_meaning:NnF #1#2
11238 { \exp:w \_token_case:NNnTF \token_if_eq_meaning:NNTF #1 {#2} { } }
11239 \cs_new:Npn \token_case_meaning:NnTF
11240 { \exp:w \_token_case:NNnTF \token_if_eq_meaning:NNTF }
11241 \cs_new:Npn \_token_case:NNnTF #1#2#3#4#5
11242 {
11243   \_token_case:NNw #1 #2 #3 #2 { }
11244   \s__token_mark {#4}
11245   \s__token_mark {#5}
11246   \s__token_stop
11247 }
11248 \cs_new:Npn \_token_case:NNw #1#2#3#4
11249 {
11250   #1 #2 #3
11251   { \_token_case_end:nw {#4} }
11252   { \_token_case:NNw #1 #2 }
11253 }

```

To tidy up the recursion, there are two outcomes. If there was a hit to one of the cases searched for, then #1 is the code to insert, #2 is the *next* case to check on and #3 is all of the rest of the cases code. That means that #4 is the **true** branch code, and #5 tidies up the spare `\s__token_mark` and the **false** branch. On the other hand, if none of the cases matched then we arrive here using the “termination” case of comparing the search with itself. That means that #1 is empty, #2 is the first `\s__token_mark` and so #4 is the **false** code (the **true** code is mopped up by #3).

```

11254 \cs_new:Npn \_token_case_end:nw #1#2#3 \s__token_mark #4#5 \s__token_stop
11255 { \exp_end: #1 #4 }

```

(End definition for `\token_case_catcode:NnTF` and others. These functions are documented on page 138.)

17.6 Peeking ahead at the next token

```

11256 (@@=peek)

```

Peeking ahead is implemented using a two part mechanism. The outer level provides a defined interface to the lower level material. This allows a large amount of code to be shared. There are four cases:

1. peek at the next token;

2. peek at the next non-space token;
3. peek at the next token and remove it;
4. peek at the next non-space token and remove it.

\l_peek_token Storage tokens which are publicly documented: the token peeked.

\g_peek_token 11257 \cs_new_eq:NN \l_peek_token ?
11258 \cs_new_eq:NN \g_peek_token ?

(End definition for \l_peek_token and \g_peek_token. These variables are documented on page 139.)

\l__peek_search_token The token to search for as an implicit token: cf. **\l__peek_search_tl**.

11259 \cs_new_eq:NN \l__peek_search_token ?

(End definition for \l__peek_search_token.)

\l__peek_search_tl The token to search for as an explicit token: cf. **\l__peek_search_token**.

11260 \tl_new:N \l__peek_search_tl

(End definition for \l__peek_search_tl.)

__peek_true:w Functions used by the branching and space-stripping code.

__peek_true_aux:w 11261 \cs_new:Npn __peek_true:w { }
__peek_false:w 11262 \cs_new:Npn __peek_true_aux:w { }
__peek_tmp:w 11263 \cs_new:Npn __peek_false:w { }
11264 \cs_new:Npn __peek_tmp:w { }

(End definition for __peek_true:w and others.)

\s__peek_mark Internal scan marks.

\s__peek_stop 11265 \scan_new:N \s__peek_mark
11266 \scan_new:N \s__peek_stop

(End definition for \s__peek_mark and \s__peek_stop.)

__peek_use_none_delimit_by_s_stop:w Functions to gobble up to a scan mark.

11267 \cs_new:Npn __peek_use_none_delimit_by_s_stop:w #1 \s__peek_stop { }

(End definition for __peek_use_none_delimit_by_s_stop:w.)

\peek_after:Nw Simple wrappers for **\futurelet**: no arguments absorbed here.

\peek_gafter:Nw 11268 \cs_new_protected:Npn \peek_after:Nw
11269 { \tex_futurelet:D \l_peek_token }
11270 \cs_new_protected:Npn \peek_gafter:Nw
11271 { \tex_global:D \tex_futurelet:D \g_peek_token }

(End definition for \peek_after:Nw and \peek_gafter:Nw. These functions are documented on page 138.)

__peek_true_remove:w A function to remove the next token and then regain control.

11272 \cs_new_protected:Npn __peek_true_remove:w
11273 {
11274 \tex_afterassignment:D __peek_true_aux:w
11275 \cs_set_eq:NN __peek_tmp:w
11276 }

(End definition for _peek_true_remove:w.)

\peek_remove_spaces:n Repeatedly use _peek_true_remove:w to remove a space and call _peek_true_remove_spaces:w.

```

11277 \cs_new_protected:Npn \peek_remove_spaces:n #1
11278 {
11279   \cs_set:Npx \_peek_false:w { \exp_not:n {#1} }
11280   \group_align_safe_begin:
11281   \cs_set:Npn \_peek_true_aux:w { \peek_after:Nw \_peek_remove_spaces: }
11282   \_peek_true_aux:w
11283 }
11284 \cs_new_protected:Npn \_peek_remove_spaces:
11285 {
11286   \if_meaning:w \l_peek_token \c_space_token
11287     \exp_after:wN \_peek_true_remove:w
11288   \else:
11289     \group_align_safe_end:
11290     \exp_after:wN \_peek_false:w
11291   \fi:
11292 }

```

(End definition for \peek_remove_spaces:n and _peek_remove_spaces:. This function is documented on page 277.)

_peek_token_generic_aux:NNTF

The generic functions store the test token in both implicit and explicit modes, and the true and false code as token lists, more or less. The two branches have to be absorbed here as the input stream needs to be cleared for the peek function itself. Here, #1 is _peek_true_remove:w when removing the token and _peek_true_aux:w otherwise.

```

11293 \cs_new_protected:Npn \_peek_token_generic_aux:NNTF #1#2#3#4#5
11294 {
11295   \group_align_safe_begin:
11296   \cs_set_eq:NN \l_peek_search_token #3
11297   \tl_set:Nn \l_peek_search_tl {#3}
11298   \cs_set:Npx \_peek_true_aux:w
11299   {
11300     \exp_not:N \group_align_safe_end:
11301     \exp_not:n {#4}
11302   }
11303   \cs_set_eq:NN \_peek_true:w #1
11304   \cs_set:Npx \_peek_false:w
11305   {
11306     \exp_not:N \group_align_safe_end:
11307     \exp_not:n {#5}
11308   }
11309   \peek_after:Nw #2
11310 }

```

(End definition for _peek_token_generic_aux:NNTF.)

_peek_token_generic:NNTF

For token removal there needs to be a call to the auxiliary function which does the work.

_peek_token_remove_generic:NNTF

```

11311 \cs_new_protected:Npn \_peek_token_generic:NNTF
11312 { \_peek_token_generic_aux:NNTF \_peek_true_aux:w }
11313 \cs_new_protected:Npn \_peek_token_generic:NNT #1#2#3
11314 { \_peek_token_generic:NNTF #1 #2 {#3} { } }

```

```

11315 \cs_new_protected:Npn \__peek_token_generic:NNTF #1#2#3
11316 { \__peek_token_generic:NNTF #1 #2 { } {#3} }
11317 \cs_new_protected:Npn \__peek_token_remove_generic:NNTF
11318 { \__peek_token_generic_aux:NNNTF \__peek_true_remove:w }
11319 \cs_new_protected:Npn \__peek_token_remove_generic:NNT #1#2#3
11320 { \__peek_token_remove_generic:NNTF #1 #2 {#3} { } }
11321 \cs_new_protected:Npn \__peek_token_remove_generic:NNTF #1#2#3
11322 { \__peek_token_remove_generic:NNTF #1 #2 { } {#3} }

```

(End definition for __peek_token_generic:NNTF and __peek_token_remove_generic:NNTF.)

__peek_execute_branches_meaning: The meaning test is straight forward.

```

11323 \cs_new:Npn \__peek_execute_branches_meaning:
11324 {
11325   \if_meaning:w \l_peek_token \l__peek_search_token
11326   \exp_after:wN \__peek_true:w
11327   \else:
11328   \exp_after:wN \__peek_false:w
11329   \fi:
11330 }

```

(End definition for __peek_execute_branches_meaning:.)

__peek_execute_branches_catcode: The catcode and charcode tests are very similar, and in order to use the same auxiliaries we do something a little bit odd, firing \if_catcode:w and \if_charcode:w before finding the operands for those tests, which are only given in the auxii:N and auxiii: auxiliaries. For our purposes, three kinds of tokens may follow the peeking function:

- control sequences which are not equal to a non-active character token (*e.g.*, macro, primitive);
- active characters which are not equal to a non-active character token (*e.g.*, macro, primitive);
- explicit non-active character tokens, or control sequences or active characters set equal to a non-active character token.

The first two cases are not distinguishable simply using TeX's \futurelet, because we can only access the \meaning of tokens in that way. In those cases, detected thanks to a comparison with \scan_stop:, we grab the following token, and compare it explicitly with the explicit search token stored in \l__peek_search_tl. The \exp_not:N prevents outer macros (coming from non-L^AT_EX3 code) from blowing up. In the third case, \l__peek_token is good enough for the test, and we compare it again with the explicit search token. Just like the peek token, the search token may be of any of the three types above, hence the need to use the explicit token that was given to the peek function.

```

11331 \cs_new:Npn \__peek_execute_branches_catcode:
11332 { \if_catcode:w \__peek_execute_branches_catcode_aux: }
11333 \cs_new:Npn \__peek_execute_branches_charcode:
11334 { \if_charcode:w \__peek_execute_branches_catcode_aux: }
11335 \cs_new:Npn \__peek_execute_branches_catcode_aux:
11336 {
11337   \if_catcode:w \exp_not:N \l_peek_token \scan_stop:
11338   \exp_after:wN \exp_after:wN
11339   \exp_after:wN \__peek_execute_branches_catcode_auxii:N

```

```

11340         \exp_after:wN \exp_not:N
11341     \else:
11342         \exp_after:wN \__peek_execute_branches_catcode_auxiii:
11343     \fi:
11344 }
11345 \cs_new:Npn \__peek_execute_branches_catcode_auxii:N #1
11346 {
11347     \exp_not:N #1
11348     \exp_after:wN \exp_not:N \l__peek_search_tl
11349     \exp_after:wN \__peek_true:w
11350 \else:
11351     \exp_after:wN \__peek_false:w
11352 \fi:
11353 #1
11354 }
11355 \cs_new:Npn \__peek_execute_branches_catcode_auxiii:
11356 {
11357     \exp_not:N \l__peek_token
11358     \exp_after:wN \exp_not:N \l__peek_search_tl
11359     \exp_after:wN \__peek_true:w
11360 \else:
11361     \exp_after:wN \__peek_false:w
11362 \fi:
11363 }

```

(End definition for __peek_execute_branches_catcode: and others.)

\peek_catcode:NTF The public functions themselves cannot be defined using \prg_new_conditional:Npnn. Instead, the TF, T, F variants are defined in terms of corresponding variants of __peek_token_generic:NNTF or __peek_token_remove_generic:NNTF, with first argument one of __peek_execute_branches_catcode:, __peek_execute_branches_charcode:, or __peek_execute_branches_meaning:.

\peek_catcode_remove:NTF

\peek_charcode:NTF

\peek_charcode_remove:NTF

\peek_meaning:NTF

\peek_meaning_remove:NTF

```

11364 \tl_map_inline:nn { { catcode } { charcode } { meaning } }
11365 {
11366     \tl_map_inline:nn { { } { _remove } }
11367     {
11368         \tl_map_inline:nn { { TF } { T } { F } }
11369         {
11370             \cs_new_protected:cpx { peek_ #1 ##1 :N ####1 }
11371             {
11372                 \exp_not:c { __peek_token ##1 _generic:NN ####1 }
11373                 \exp_not:c { __peek_execute_branches_ #1 : }
11374             }
11375         }
11376     }
11377 }

```

(End definition for \peek_catcode:NTF and others. These functions are documented on page 139.)

\peek_catcode_ignore_spaces:NTF To ignore spaces, remove them using \peek_remove_spaces:n before running the tests.

\peek_catcode_remove_ignore_spaces:NTF

\peek_charcode_ignore_spaces:NTF

\peek_charcode_remove_ignore_spaces:NTF

\peek_meaning_ignore_spaces:NTF

\peek_meaning_remove_ignore_spaces:NTF

```

11378 \tl_map_inline:nn
11379 {
11380     { catcode } { catcode_remove }
11381     { charcode } { charcode_remove }

```

```

11382 { meaning } { meaning_remove }
11383 }
11384 {
11385   \cs_new_protected:cpx { peek_#1_ignore_spaces:NTF } ##1##2##3
11386   {
11387     \peek_remove_spaces:n
11388     { \exp_not:c { peek_#1:NTF } ##1 {##2} {##3} }
11389   }
11390   \cs_new_protected:cpx { peek_#1_ignore_spaces:NT } ##1##2
11391   {
11392     \peek_remove_spaces:n
11393     { \exp_not:c { peek_#1:NT } ##1 {##2} }
11394   }
11395   \cs_new_protected:cpx { peek_#1_ignore_spaces:NF } ##1##2
11396   {
11397     \peek_remove_spaces:n
11398     { \exp_not:c { peek_#1:NF } ##1 {##2} }
11399   }
11400 }

```

(End definition for `\peek_catcode_ignore_spaces:NTF` and others. These functions are documented on page 139.)

\peek_N_type:TF All tokens are N-type tokens, except in four cases: begin-group tokens, end-group tokens, space tokens with character code 32, and outer tokens. Since `\l_peek_token` might be outer, we cannot use the convenient `\bool_if:nTF` function, and must resort to the old trick of using `\ifodd` to expand a set of tests. The **false** branch of this test is taken if the token is one of the first three kinds of non-N-type tokens (explicit or implicit), thus we call `__peek_false:w`. In the **true** branch, we must detect outer tokens, without impacting performance too much for non-outer tokens. The first filter is to search for **outer** in the `\meaning` of `\l_peek_token`. If that is absent, `__peek_use_none_delimit_by_s_stop:w` cleans up, and we call `__peek_true:w`. Otherwise, the token can be a non-outer macro or a primitive mark whose parameter or replacement text contains **outer**, it can be the primitive `\outer`, or it can be an outer token. Macros and marks would have **ma** in the part before the first occurrence of **outer**; the meaning of `\outer` has nothing after **outer**, contrarily to outer macros; and that covers all cases, calling `__peek_true:w` or `__peek_false:w` as appropriate. Here, there is no *search token*, so we feed a dummy `\scan_stop:` to the `__peek_token_generic:NNTF` function.

```

11401 \group_begin:
11402   \cs_set_protected:Npn \__peek_tmp:w #1 \s__peek_stop
11403   {
11404     \cs_new_protected:Npn \__peek_execute_branches_N_type:
11405     {
11406       \if_int_odd:w
11407         \if_catcode:w \exp_not:N \l_peek_token { 0 \exp_stop_f: \fi:
11408         \if_catcode:w \exp_not:N \l_peek_token } 0 \exp_stop_f: \fi:
11409         \if_meaning:w \l_peek_token \c_space_token 0 \exp_stop_f: \fi:
11410         1 \exp_stop_f:
11411       \exp_after:wN \__peek_N_type:w
11412       \token_to_meaning:N \l_peek_token
11413       \s__peek_mark \__peek_N_type_aux:nnw
11414       #1 \s__peek_mark \__peek_use_none_delimit_by_s_stop:w
11415       \s__peek_stop

```

```

11416         \exp_after:wN \__peek_true:w
11417     \else:
11418         \exp_after:wN \__peek_false:w
11419     \fi:
11420 }
11421 \cs_new_protected:Npn \__peek_N_type:w ##1 #1 ##2 \s__peek_mark ##3
11422 { ##3 {##1} {##2} }
11423 }
11424 \exp_after:wN \__peek_tmp:w \tl_to_str:n { outer } \s__peek_stop
11425 \group_end:
11426 \cs_new_protected:Npn \__peek_N_type_aux:nnw #1 #2 #3 \fi:
11427 {
11428     \fi:
11429     \tl_if_in:noTF {#1} { \tl_to_str:n {ma} }
11430     { \__peek_true:w }
11431     { \tl_if_empty:nTF {#2} { \__peek_true:w } { \__peek_false:w } }
11432 }
11433 \cs_new_protected:Npn \peek_N_type:TF
11434 {
11435     \__peek_token_generic:NNTF
11436     \__peek_execute_branches_N_type: \scan_stop:
11437 }
11438 \cs_new_protected:Npn \peek_N_type:T
11439 { \__peek_token_generic:NNT \__peek_execute_branches_N_type: \scan_stop: }
11440 \cs_new_protected:Npn \peek_N_type:F
11441 { \__peek_token_generic:NNTF \__peek_execute_branches_N_type: \scan_stop: }

```

(End definition for \peek_N_type:TF and others. This function is documented on page 141.)

```

11442 \</tex>
11443 \</package>

```

18 l3prop implementation

The following test files are used for this code: `m3prop001`, `m3prop002`, `m3prop003`, `m3prop004`, `m3show001`.

```

11444 \*package>
11445 \@@=prop>

```

A property list is a macro whose top-level expansion is of the form

```

\s__prop \__prop_pair:wn <key1> \s__prop {<value1>}
...
\__prop_pair:wn <keyn> \s__prop {<valuen>}

```

where `\s__prop` is a scan mark (equal to `\scan_stop:`), and `__prop_pair:wn` can be used to map through the property list.

`\s__prop` The internal token used at the beginning of property lists. This is also used after each `<key>` (see `__prop_pair:wn`).

(End definition for `\s__prop`.)

<code>__prop_pair:wn</code>	<p><code>__prop_pair:wn <key> \s__prop {(item)}</code></p> <p>The internal token used to begin each key–value pair in the property list. If expanded outside of a mapping or manipulation function, an error is raised. The definition should always be set globally.</p> <p><i>(End definition for __prop_pair:wn.)</i></p>
<code>\l__prop_internal_tl</code>	<p>Token list used to store new key–value pairs to be inserted by functions of the <code>\prop_put:Nnn</code> family.</p> <p><i>(End definition for \l__prop_internal_tl.)</i></p>
<hr/> <code>__prop_split:NnTF</code> <hr/> Updated: 2013-01-08 <hr/>	<p><code>__prop_split:NnTF <property list> {(key)} {(true code)} {(false code)}</code></p> <p>Splits the <i><property list></i> at the <i><key></i>, giving three token lists: the <i><extract></i> of <i><property list></i> before the <i><key></i>, the <i><value></i> associated with the <i><key></i> and the <i><extract></i> of the <i><property list></i> after the <i><value></i>. Both <i><extracts></i> retain the internal structure of a property list, and the concatenation of the two <i><extracts></i> is a property list. If the <i><key></i> is present in the <i><property list></i> then the <i><true code></i> is left in the input stream, with #1, #2, and #3 replaced by the first <i><extract></i>, the <i><value></i>, and the second <i><extract></i>. If the <i><key></i> is not present in the <i><property list></i> then the <i><false code></i> is left in the input stream, with no trailing material. Both <i><true code></i> and <i><false code></i> are used in the replacement text of a macro defined internally, hence macro parameter characters should be doubled, except #1, #2, and #3 which stand in the <i><true code></i> for the three extracts from the property list. The <i><key></i> comparison takes place as described for <code>\str_if_eq:nn</code>.</p>
<code>\s__prop</code>	<p>A private scan mark is used as a marker after each key, and at the very beginning of the property list.</p> <p>11446 <code>\scan_new:N \s__prop</code></p> <p><i>(End definition for \s__prop.)</i></p>
<code>__prop_pair:wn</code>	<p>The delimiter is always defined, but when misused simply triggers an error and removes its argument.</p> <p>11447 <code>\cs_new:Npn __prop_pair:wn #1 \s__prop #2</code> 11448 <code>{ __kernel_msg_expandable_error:nn { kernel } { misused-prop } }</code></p> <p><i>(End definition for __prop_pair:wn.)</i></p>
<code>\l__prop_internal_tl</code>	<p>Token list used to store the new key–value pair inserted by <code>\prop_put:Nnn</code> and friends.</p> <p>11449 <code>\tl_new:N \l__prop_internal_tl</code></p> <p><i>(End definition for \l__prop_internal_tl.)</i></p>
<code>\c_empty_prop</code>	<p>An empty prop.</p> <p>11450 <code>\tl_const:Nn \c_empty_prop { \s__prop }</code></p> <p><i>(End definition for \c_empty_prop. This variable is documented on page 152.)</i></p>

18.1 Internal auxiliaries

`\s__prop_mark` Internal scan marks.

`\s__prop_stop` 11451 `\scan_new:N \s__prop_mark`
11452 `\scan_new:N \s__prop_stop`

(End definition for `\s__prop_mark` and `\s__prop_stop`.)

`\q__prop_recursion_tail` Internal recursion quarks.

`\q__prop_recursion_stop` 11453 `\quark_new:N \q__prop_recursion_tail`
11454 `\quark_new:N \q__prop_recursion_stop`

(End definition for `\q__prop_recursion_tail` and `\q__prop_recursion_stop`.)

`__prop_if_recursion_tail_stop:n` Functions to query recursion quarks.

`__prop_if_recursion_tail_stop:o` 11455 `__kernel_quark_new_test:N __prop_if_recursion_tail_stop:n`
11456 `\cs_generate_variant:Nn __prop_if_recursion_tail_stop:n { o }`

(End definition for `__prop_if_recursion_tail_stop:n` and `__prop_if_recursion_tail_stop:o`.)

18.2 Allocation and initialisation

`\prop_new:N` Property lists are initialized with the value `\c_empty_prop`.

`\prop_new:c` 11457 `\cs_new_protected:Npn \prop_new:N #1`
11458 `{`
11459 `__kernel_chk_if_free_cs:N #1`
11460 `\cs_gset_eq:NN #1 \c_empty_prop`
11461 `}`
11462 `\cs_generate_variant:Nn \prop_new:N { c }`

(End definition for `\prop_new:N`. This function is documented on page 146.)

`\prop_clear:N` The same idea for clearing.

`\prop_clear:c` 11463 `\cs_new_protected:Npn \prop_clear:N #1`
`\prop_gclear:N` 11464 `{ \prop_set_eq:NN #1 \c_empty_prop }`
`\prop_gclear:c` 11465 `\cs_generate_variant:Nn \prop_clear:N { c }`
11466 `\cs_new_protected:Npn \prop_gclear:N #1`
11467 `{ \prop_gset_eq:NN #1 \c_empty_prop }`
11468 `\cs_generate_variant:Nn \prop_gclear:N { c }`

(End definition for `\prop_clear:N` and `\prop_gclear:N`. These functions are documented on page 146.)

`\prop_clear_new:N` Once again a simple variation of the token list functions.

`\prop_clear_new:c` 11469 `\cs_new_protected:Npn \prop_clear_new:N #1`
`\prop_gclear_new:N` 11470 `{ \prop_if_exist:NTF #1 { \prop_clear:N #1 } { \prop_new:N #1 } }`
`\prop_gclear_new:c` 11471 `\cs_generate_variant:Nn \prop_clear_new:N { c }`
11472 `\cs_new_protected:Npn \prop_gclear_new:N #1`
11473 `{ \prop_if_exist:NTF #1 { \prop_gclear:N #1 } { \prop_new:N #1 } }`
11474 `\cs_generate_variant:Nn \prop_gclear_new:N { c }`

(End definition for `\prop_clear_new:N` and `\prop_gclear_new:N`. These functions are documented on page 146.)

```

\prop_set_eq:NN These are simply copies from the token list functions.
\prop_set_eq:cN 11475 \cs_new_eq:NN \prop_set_eq:NN \tl_set_eq:NN
\prop_set_eq:Nc 11476 \cs_new_eq:NN \prop_set_eq:Nc \tl_set_eq:Nc
\prop_set_eq:cc 11477 \cs_new_eq:NN \prop_set_eq:cN \tl_set_eq:cN
\prop_gset_eq:NN 11478 \cs_new_eq:NN \prop_set_eq:cc \tl_set_eq:cc
\prop_gset_eq:cN 11479 \cs_new_eq:NN \prop_gset_eq:NN \tl_gset_eq:NN
\prop_gset_eq:Nc 11480 \cs_new_eq:NN \prop_gset_eq:Nc \tl_gset_eq:Nc
\prop_gset_eq:cN 11481 \cs_new_eq:NN \prop_gset_eq:cN \tl_gset_eq:cN
\prop_gset_eq:cc 11482 \cs_new_eq:NN \prop_gset_eq:cc \tl_gset_eq:cc

```

(End definition for `\prop_set_eq:NN` and `\prop_gset_eq:NN`. These functions are documented on page 146.)

```

\l_tmpa_prop We can now initialize the scratch variables.
\l_tmpb_prop
\g_tmpa_prop 11483 \prop_new:N \l_tmpa_prop
\g_tmpb_prop 11484 \prop_new:N \l_tmpb_prop
\g_tmpb_prop 11485 \prop_new:N \g_tmpa_prop
\g_tmpb_prop 11486 \prop_new:N \g_tmpb_prop

```

(End definition for `\l_tmpa_prop` and others. These variables are documented on page 151.)

```

\l__prop_internal_prop Property list used by \prop_set_from_keyval:Nn and others.
11487 \prop_new:N \l__prop_internal_prop

```

(End definition for `\l__prop_internal_prop`.)

```

\prop_set_from_keyval:Nn To avoid tracking throughout the loop the variable name and whether the assignment
\prop_set_from_keyval:cN is local/global, do everything in a scratch variable and empty it afterwards to avoid
\prop_gset_from_keyval:Nn wasting memory. Loop through items separated by commas, with \prg_do_nothing:
\prop_gset_from_keyval:cN to avoid losing braces. After checking for termination, split the item at the first and
\prop_const_from_keyval:Nn then at the second = (which ought to be the first of the trailing = that we added). For
\prop_const_from_keyval:cN both splits trim spaces and call a function (first \__prop_from_keyval_key:w then \__-
\__prop_from_keyval:n prop_from_keyval_value:w), followed by the trimmed material, \s__prop_mark, the
\__prop_from_keyval_loop:w subsequent part of the item, and the trailing ='s and \s__prop_stop. After finding
\__prop_from_keyval_split:Nw the <key> just store it after \s__prop_stop. After finding the <value> ignore completely
\__prop_from_keyval_key:n empty items (both trailing = were used as delimiters and all parts are empty); if the
\__prop_from_keyval_key:w remaining part #2 consists exactly of the second trailing = (namely there was exactly one
\__prop_from_keyval_value:n = in the item) then output one key–value pair for the property list; otherwise complain
\__prop_from_keyval_value:w about a missing or extra =.

```

```

11488 \cs_new_protected:Npn \prop_set_from_keyval:Nn #1#2
11489 {
11490   \prop_clear:N \l__prop_internal_prop
11491   \__prop_from_keyval:n {#2}
11492   \prop_set_eq:NN #1 \l__prop_internal_prop
11493   \prop_clear:N \l__prop_internal_prop
11494 }
11495 \cs_generate_variant:Nn \prop_set_from_keyval:Nn { c }
11496 \cs_new_protected:Npn \prop_gset_from_keyval:Nn #1#2
11497 {
11498   \prop_clear:N \l__prop_internal_prop
11499   \__prop_from_keyval:n {#2}
11500   \prop_gset_eq:NN #1 \l__prop_internal_prop
11501   \prop_clear:N \l__prop_internal_prop

```

```

11502 }
11503 \cs_generate_variant:Nn \prop_gset_from_keyval:Nn { c }
11504 \cs_new_protected:Npn \prop_const_from_keyval:Nn #1#2
11505 {
11506   \prop_clear:N \l__prop_internal_prop
11507   \__prop_from_keyval:n {#2}
11508   \tl_const:Nx #1 { \exp_not:o \l__prop_internal_prop }
11509   \prop_clear:N \l__prop_internal_prop
11510 }
11511 \cs_generate_variant:Nn \prop_const_from_keyval:Nn { c }
11512 \cs_new_protected:Npn \__prop_from_keyval:n #1
11513 {
11514   \__prop_from_keyval_loop:w \prg_do_nothing: #1 ,
11515   \q__prop_recursion_tail , \q__prop_recursion_stop
11516 }
11517 \cs_new_protected:Npn \__prop_from_keyval_loop:w #1 ,
11518 {
11519   \__prop_if_recursion_tail_stop:o {#1}
11520   \__prop_from_keyval_split:Nw \__prop_from_keyval_key:n
11521   #1 = = \s__prop_stop {#1}
11522   \__prop_from_keyval_loop:w \prg_do_nothing:
11523 }
11524 \cs_new_protected:Npn \__prop_from_keyval_split:Nw #1#2 =
11525 { \tl_trim_spaces_apply:oN {#2} #1 }
11526 \cs_new_protected:Npn \__prop_from_keyval_key:n #1
11527 { \__prop_from_keyval_key:w #1 \s__prop_mark }
11528 \cs_new_protected:Npn \__prop_from_keyval_key:w #1 \s__prop_mark #2 \s__prop_stop
11529 {
11530   \__prop_from_keyval_split:Nw \__prop_from_keyval_value:n
11531   \prg_do_nothing: #2 \s__prop_stop {#1}
11532 }
11533 \cs_new_protected:Npn \__prop_from_keyval_value:n #1
11534 { \__prop_from_keyval_value:w #1 \s__prop_mark }
11535 \cs_new_protected:Npn \__prop_from_keyval_value:w #1 \s__prop_mark #2 \s__prop_stop #3#4
11536 {
11537   \tl_if_empty:nF { #3 #1 #2 }
11538   {
11539     \str_if_eq:nnTF {#2} { = }
11540     { \prop_put:Nnn \l__prop_internal_prop {#3} {#1} }
11541     {
11542       \__kernel_msg_error:nnx { kernel } { prop-keyval }
11543       { \exp_not:o {#4} }
11544     }
11545   }
11546 }

```

(End definition for `\prop_set_from_keyval:Nn` and others. These functions are documented on page 146.)

18.3 Accessing data in property lists

`__prop_split:NnTF` This function is used by most of the module, and hence must be fast. It receives a *⟨property list⟩*, a *⟨key⟩*, a *⟨true code⟩* and a *⟨false code⟩*. The aim is to split the *⟨property list⟩* at the given *⟨key⟩* into the *⟨extract₁⟩* before the key–value pair, the *⟨value⟩* associated

with the $\langle key \rangle$ and the $\langle extract_2 \rangle$ after the key-value pair. This is done using a delimited function, whose definition is as follows, where the $\langle key \rangle$ is turned into a string.

```
\cs_set:Npn \__prop_split_aux:w #1
\__prop_pair:wn \langle key \rangle \s__prop #2
#3 \s__prop_mark #4 #5 \s__prop_stop
{ #4 { \langle true code \rangle } { \langle false code \rangle } }
```

If the $\langle key \rangle$ is present in the property list, $\backslash__prop_split_aux:w$'s #1 is the part before the $\langle key \rangle$, #2 is the $\langle value \rangle$, #3 is the part after the $\langle key \rangle$, #4 is $\backslash use_i:nn$, and #5 is additional tokens that we do not care about. The $\langle true code \rangle$ is left in the input stream, and can use the parameters #1, #2, #3 for the three parts of the property list as desired. Namely, the original property list is in this case #1 $\backslash__prop_pair:wn \langle key \rangle \backslash s_prop \{ \#2 \}$ #3.

If the $\langle key \rangle$ is not there, then the $\langle function \rangle$ is $\backslash use_ii:nn$, which keeps the $\langle false code \rangle$.

```
11547 \cs_new_protected:Npn \__prop_split:NnTF #1#2
11548 { \exp_args:NNo \__prop_split_aux:NnTF #1 { \tl_to_str:n {#2} } }
11549 \cs_new_protected:Npn \__prop_split_aux:NnTF #1#2#3#4
11550 {
11551   \cs_set:Npn \__prop_split_aux:w ##1
11552     \__prop_pair:wn #2 \s__prop ##2 ##3 \s__prop_mark ##4 ##5 \s__prop_stop
11553     { ##4 {#3} {#4} }
11554   \exp_after:wN \__prop_split_aux:w #1 \s__prop_mark \use_i:nn
11555   \__prop_pair:wn #2 \s__prop { } \s__prop_mark \use_ii:nn \s__prop_stop
11556 }
11557 \cs_new:Npn \__prop_split_aux:w { }
```

(End definition for $\backslash__prop_split:NnTF$, $\backslash__prop_split_aux:NnTF$, and $\backslash__prop_split_aux:w$.)

$\backslash prop_remove:Nn$ $\backslash prop_remove:NV$ $\backslash prop_remove:cn$ $\backslash prop_remove:cV$ $\backslash prop_gremove:Nn$ $\backslash prop_gremove:NV$ $\backslash prop_gremove:cn$ $\backslash prop_gremove:cV$	Deleting from a property starts by splitting the list. If the key is present in the property list, the returned value is ignored. If the key is missing, nothing happens. <pre>11558 \cs_new_protected:Npn \prop_remove:Nn #1#2 11559 { 11560 __prop_split:NnTF #1 {#2} 11561 { \tl_set:Nn #1 { ##1 ##3 } } 11562 { } 11563 } 11564 \cs_new_protected:Npn \prop_gremove:Nn #1#2 11565 { 11566 __prop_split:NnTF #1 {#2} 11567 { \tl_gset:Nn #1 { ##1 ##3 } } 11568 { } 11569 } 11570 \cs_generate_variant:Nn \prop_remove:Nn { NV } 11571 \cs_generate_variant:Nn \prop_remove:Nn { c , cV } 11572 \cs_generate_variant:Nn \prop_gremove:Nn { NV } 11573 \cs_generate_variant:Nn \prop_gremove:Nn { c , cV }</pre>
--	---

(End definition for $\backslash prop_remove:Nn$ and $\backslash prop_gremove:Nn$. These functions are documented on page 148.)

\prop_get:NnN Getting an item from a list is very easy: after splitting, if the key is in the property list, just set the token list variable to the return value, otherwise to `\q_no_value`.

```

\prop_get:NVN
\prop_get:NoN 11574 \cs_new_protected:Npn \prop_get:NnN #1#2#3
\prop_get:cnN 11575 {
\prop_get:cVN 11576   \__prop_split:NnTF #1 {#2}
\prop_get:coN 11577   { \tl_set:Nn #3 {##2} }
                { \tl_set:Nn #3 { \q_no_value } }
11579 }
11580 \cs_generate_variant:Nn \prop_get:NnN { NV , Nv , No }
11581 \cs_generate_variant:Nn \prop_get:NnN { c , cV , cv , co }
```

(End definition for `\prop_get:NnN`. This function is documented on page 147.)

\prop_pop:NnN Popping a value also starts by doing the split. If the key is present, save the value in the token list and update the property list as when deleting. If the key is missing, save `\q_no_value` in the token list.

```

\prop_pop:NoN 11582 \cs_new_protected:Npn \prop_pop:NnN #1#2#3
\prop_pop:cnN 11583 {
\prop_pop:coN 11584   \__prop_split:NnTF #1 {#2}
\prop_gpop:NnN 11585   {
\prop_gpop:NoN 11586     \tl_set:Nn #3 {##2}
\prop_gpop:cnN 11587     \tl_set:Nn #1 { ##1 ##3 }
11588   }
11589   { \tl_set:Nn #3 { \q_no_value } }
11590 }
11591 \cs_new_protected:Npn \prop_gpop:NnN #1#2#3
11592 {
11593   \__prop_split:NnTF #1 {#2}
11594   {
11595     \tl_set:Nn #3 {##2}
11596     \tl_gset:Nn #1 { ##1 ##3 }
11597   }
11598   { \tl_set:Nn #3 { \q_no_value } }
11599 }
11600 \cs_generate_variant:Nn \prop_pop:NnN { No }
11601 \cs_generate_variant:Nn \prop_pop:NnN { c , co }
11602 \cs_generate_variant:Nn \prop_gpop:NnN { No }
11603 \cs_generate_variant:Nn \prop_gpop:NnN { c , co }
```

(End definition for `\prop_pop:NnN` and `\prop_gpop:NnN`. These functions are documented on page 147.)

\prop_item:Nn Getting the value corresponding to a key in a property list in an expandable fashion is similar to mapping some tokens. Go through the property list one $\langle key \rangle$ – $\langle value \rangle$ pair at a time: the arguments of `__prop_item_Nn:nwn` are the $\langle key \rangle$ we are looking for, a $\langle key \rangle$ of the property list, and its associated value. The $\langle keys \rangle$ are compared (as strings). If they match, the $\langle value \rangle$ is returned, within `\exp_not:n`. The loop terminates even if the $\langle key \rangle$ is missing, and yields an empty value, because we have appended the appropriate $\langle key \rangle$ – $\langle empty value \rangle$ pair to the property list.

```

11604 \cs_new:Npn \prop_item:Nn #1#2
11605 {
11606   \exp_last_unbraced:Noo \__prop_item_Nn:nwn { \tl_to_str:n {#2} } #1
11607   \__prop_pair:wn \tl_to_str:n {#2} \s_prop { }
11608   \prg_break_point:
```

```

11609 }
11610 \cs_new:Npn \__prop_item:Nn:nwwn #1#2 \__prop_pair:wn #3 \s__prop #4
11611 {
11612   \str_if_eq:eeTF {#1} {#3}
11613   { \prg_break:n { \exp_not:n {#4} } }
11614   { \__prop_item:Nn:nwwn {#1} }
11615 }
11616 \cs_generate_variant:Nn \prop_item:Nn { c }

```

(End definition for `\prop_item:Nn` and `__prop_item:Nn:nwwn`. This function is documented on page 148.)

`\prop_count:N` Counting the key–value pairs in a property list is done using the same approach as for
`\prop_count:c` other count functions: turn each entry into a +1 then use integer evaluation to actually
`__prop_count:nn` do the mathematics.

```

11617 \cs_new:Npn \prop_count:N #1
11618 {
11619   \int_eval:n
11620   {
11621     0
11622     \prop_map_function:NN #1 \__prop_count:nn
11623   }
11624 }
11625 \cs_new:Npn \__prop_count:nn #1#2 { + 1 }
11626 \cs_generate_variant:Nn \prop_count:N { c }

```

(End definition for `\prop_count:N` and `__prop_count:nn`. This function is documented on page 148.)

`\prop_pop:NnNTF` Popping an item from a property list, keeping track of whether the key was present or
`\prop_pop:cnNTF` not, is implemented as a conditional. If the key was missing, neither the property list, nor
`\prop_gpop:NnNTF` the token list are altered. Otherwise, `\prg_return_true:` is used after the assignments.
`\prop_gpop:cnNTF`

```

11627 \prg_new_protected_conditional:Npnn \prop_pop:NnN #1#2#3 { T , F , TF }
11628 {
11629   \__prop_split:NnTF #1 {#2}
11630   {
11631     \tl_set:Nn #3 {##2}
11632     \tl_set:Nn #1 { ##1 ##3 }
11633     \prg_return_true:
11634   }
11635   { \prg_return_false: }
11636 }
11637 \prg_new_protected_conditional:Npnn \prop_gpop:NnN #1#2#3 { T , F , TF }
11638 {
11639   \__prop_split:NnTF #1 {#2}
11640   {
11641     \tl_set:Nn #3 {##2}
11642     \tl_gset:Nn #1 { ##1 ##3 }
11643     \prg_return_true:
11644   }
11645   { \prg_return_false: }
11646 }
11647 \prg_generate_conditional_variant:Nnn \prop_pop:NnN { c } { T , F , TF }
11648 \prg_generate_conditional_variant:Nnn \prop_gpop:NnN { c } { T , F , TF }

```

(End definition for `\prop_pop:NnTF` and `\prop_gpop:NnTF`. These functions are documented on page 149.)

`\prop_put:Nnn` Since the branches of `__prop_split:NnTF` are used as the replacement text of an internal macro, and since the $\langle key \rangle$ and new $\langle value \rangle$ may contain arbitrary tokens, it is not safe to include them in the argument of `__prop_split:NnTF`. We thus start by storing in `\l__prop_internal_tl` tokens which (after x-expansion) encode the key–value pair. This variable can safely be used in `__prop_split:NnTF`. If the $\langle key \rangle$ was absent, append the new key–value to the list. Otherwise concatenate the extracts `##1` and `##3` with the new key–value pair `\l__prop_internal_tl`. The updated entry is placed at the same spot as the original $\langle key \rangle$ in the property list, preserving the order of entries.

```

11649 \cs_new_protected:Npn \prop_put:Nnn { \__prop_put:NNnn \__kernel_tl_set:Nx }
11650 \cs_new_protected:Npn \prop_gput:Nnn { \__prop_put:NNnn \__kernel_tl_gset:Nx }
11651 \cs_new_protected:Npn \__prop_put:NNnn #1#2#3#4
11652 {
11653   \tl_set:Nn \l__prop_internal_tl
11654   {
11655     \exp_not:N \__prop_pair:wn \tl_to_str:n {#3}
11656     \s__prop { \exp_not:n {#4} }
11657   }
11658   \__prop_split:NnTF #2 {#3}
11659   { #1 #2 { \exp_not:n {##1} \l__prop_internal_tl \exp_not:n {##3} } }
11660   { #1 #2 { \exp_not:o {#2} \l__prop_internal_tl } }
11661 }
11662 \cs_generate_variant:Nn \prop_put:Nnn
11663 { NnV , Nno , Nnx , NV , NVV , NVx , Nvx , No , Noo , Nxx }
11664 \cs_generate_variant:Nn \prop_gput:Nnn
11665 { c , cnV , cno , cnx , cV , cVV , cVx , cvx , co , coo , cxx }
11666 \cs_generate_variant:Nn \prop_gput:Nnn
11667 { NnV , Nno , Nnx , NV , NVV , NVx , Nvx , No , Noo , Nxx }
11668 \cs_generate_variant:Nn \prop_gput:Nnn
11669 { c , cnV , cno , cnx , cV , cVV , cVx , cvx , co , coo , cxx }

```

(End definition for `\prop_put:Nnn` and others. These functions are documented on page 147.)

`\prop_put_if_new:Nnn` Adding conditionally also splits. If the key is already present, the three brace groups given by `__prop_split:NnTF` are removed. If the key is new, then the value is added, being careful to convert the key to a string using `\tl_to_str:n`.

```

11670 \cs_new_protected:Npn \prop_put_if_new:Nnn
11671 { \__prop_put_if_new:NNnn \__kernel_tl_set:Nx }
11672 \cs_new_protected:Npn \prop_gput_if_new:Nnn
11673 { \__prop_put_if_new:NNnn \__kernel_tl_gset:Nx }
11674 \cs_new_protected:Npn \__prop_put_if_new:NNnn #1#2#3#4
11675 {
11676   \tl_set:Nn \l__prop_internal_tl
11677   {
11678     \exp_not:N \__prop_pair:wn \tl_to_str:n {#3}
11679     \s__prop \exp_not:n { {#4} }
11680   }
11681   \__prop_split:NnTF #2 {#3}
11682   { }
11683   { #1 #2 { \exp_not:o {#2} \l__prop_internal_tl } }
11684 }

```



```

11685 \cs_generate_variant:Nn \prop_put_if_new:Nnn { c }
11686 \cs_generate_variant:Nn \prop_gput_if_new:Nnn { c }

```

(End definition for `\prop_put_if_new:Nnn`, `\prop_gput_if_new:Nnn`, and `__prop_put_if_new:NNnn`. These functions are documented on page 147.)

18.4 Property list conditionals

```

\prop_if_exist_p:N
\prop_if_exist_p:c
\prop_if_exist:NTF
\prop_if_exist:cTF

```

Copies of the `cs` functions defined in `l3basics`.

```

11687 \prg_new_eq_conditional:NNn \prop_if_exist:N \cs_if_exist:N
11688 { TF , T , F , p }
11689 \prg_new_eq_conditional:NNn \prop_if_exist:c \cs_if_exist:c
11690 { TF , T , F , p }

```

(End definition for `\prop_if_exist:NTF`. This function is documented on page 148.)

```

\prop_if_empty_p:N
\prop_if_empty_p:c
\prop_if_empty:NTF
\prop_if_empty:cTF

```

Same test as for token lists.

```

11691 \prg_new_conditional:Npnn \prop_if_empty:N #1 { p , T , F , TF }
11692 {
11693   \tl_if_eq:NNTF #1 \c_empty_prop
11694   \prg_return_true: \prg_return_false:
11695 }
11696 \prg_generate_conditional_variant:Nnn \prop_if_empty:N
11697 { c } { p , T , F , TF }

```

(End definition for `\prop_if_empty:NTF`. This function is documented on page 148.)

```

\prop_if_in_p:Nn
\prop_if_in_p:Nv
\prop_if_in_p:No
\prop_if_in_p:cn
\prop_if_in_p:cV
\prop_if_in_p:co
\prop_if_in:NnTF
\prop_if_in:NvTF
\prop_if_in:NoTF
\prop_if_in:cnTF
\prop_if_in:cVTF
\prop_if_in:coTF
\__prop_if_in:nwnn
\__prop_if_in:N

```

Testing expandably if a key is in a property list requires to go through the key–value pairs one by one. This is rather slow, and a faster test would be

```

\prg_new_protected_conditional:Npnn \prop_if_in:Nn #1 #2
{
  \@@_split:NnTF #1 {#2}
  { \prg_return_true: }
  { \prg_return_false: }
}

```

but `__prop_split:NnTF` is non-expandable.

Instead, the key is compared to each key in turn using `\str_if_eq:ee`, which is expandable. To terminate the mapping, we append to the property list the key that is searched for. This second `\tl_to_str:n` is not expanded at the start, but only when included in the `\str_if_eq:ee`. It cannot make the breaking mechanism choke, because the arbitrary token list material is enclosed in braces. The second argument of `__prop_if_in:nwnn` is most often empty. When the *key* is found in the list, `__prop_if_in:N` receives `__prop_pair:wn`, and if it is found as the extra item, the function receives `\q_prop_recursion_tail`, easily recognizable.

Here, `\prop_map_function:NN` is not sufficient for the mapping, since it can only map a single token, and cannot carry the key that is searched for.

```

11698 \prg_new_conditional:Npnn \prop_if_in:Nn #1#2 { p , T , F , TF }
11699 {
11700   \exp_last_unbraced:Noo \__prop_if_in:nwnn { \tl_to_str:n {#2} } #1
11701   \__prop_pair:wn \tl_to_str:n {#2} \s_prop { }
11702   \q_prop_recursion_tail
11703   \prg_break_point:

```

```

11704 }
11705 \cs_new:Npn \__prop_if_in:nwn #1#2 \__prop_pair:wn #3 \s__prop #4
11706 {
11707   \str_if_eq:eeTF {#1} {#3}
11708   { \__prop_if_in:N }
11709   { \__prop_if_in:nwn {#1} }
11710 }
11711 \cs_new:Npn \__prop_if_in:N #1
11712 {
11713   \if_meaning:w \q__prop_recursion_tail #1
11714   \prg_return_false:
11715   \else:
11716     \prg_return_true:
11717   \fi:
11718   \prg_break:
11719 }
11720 \prg_generate_conditional_variant:Nnn \prop_if_in:Nn
11721 { NV , No , c , cV , co } { p , T , F , TF }

```

(End definition for `\prop_if_in:NnTF`, `__prop_if_in:nwn`, and `__prop_if_in:N`. This function is documented on page 149.)

18.5 Recovering values from property lists with branching

`\prop_get:NnTF` Getting the value corresponding to a key, keeping track of whether the key was present or not, is implemented as a conditional (with side effects). If the key was absent, the token list is not altered.

```

11722 \prg_new_protected_conditional:Npnn \prop_get:NnN #1#2#3 { T , F , TF }
11723 {
11724   \__prop_split:NnTF #1 {#2}
11725   {
11726     \tl_set:Nn #3 {##2}
11727     \prg_return_true:
11728   }
11729   { \prg_return_false: }
11730 }
11731 \prg_generate_conditional_variant:Nnn \prop_get:NnN
11732 { NV , Nv , No , c , cV , cv , co } { T , F , TF }

```

(End definition for `\prop_get:NnTF`. This function is documented on page 149.)

18.6 Mapping to property lists

The argument delimited by `__prop_pair:wn` is empty except at the end of the loop where it is `\prg_break:.` No need for any quark test.

```

\prop_map_function:NN The argument delimited by \__prop_pair:wn is empty except at the end of the loop
\prop_map_function:Nc where it is \prg_break:. No need for any quark test.
\prop_map_function:cN
\prop_map_function:cc
\__prop_map_function:Nwn
11733 \cs_new:Npn \prop_map_function:NN #1#2
11734 {
11735   \exp_after:wN \use_i_ii:nnn
11736   \exp_after:wN \__prop_map_function:Nwn
11737   \exp_after:wN #2
11738   #1
11739   \prg_break: \__prop_pair:wn \s__prop { } \prg_break_point:
11740   \prg_break_point:Nn \prop_map_break: { }
11741 }

```

```

11742 \cs_new:Npn \__prop_map_function:Nwwn #1#2 \__prop_pair:wn #3 \s__prop #4
11743 {
11744     #2
11745     #1 {#3} {#4}
11746     \__prop_map_function:Nwwn #1
11747 }
11748 \cs_generate_variant:Nn \prop_map_function:NN { Nc , c , cc }

```

(End definition for `\prop_map_function:NN` and `__prop_map_function:Nwwn`. This function is documented on page 150.)

\prop_map_inline:Nn Mapping in line requires a nesting level counter. Store the current definition of `__prop_pair:wn`, and define it anew. At the end of the loop, revert to the earlier definition. Note that besides pairs of the form `__prop_pair:wn <key> \s__prop {<value>}`, there are a leading and a trailing tokens, but both are equal to `\scan_stop:`, hence have no effect in such inline mapping. Such `\scan_stop:` could have affected ligatures if they appeared during the mapping.

\prop_map_inline:cn

```

11749 \cs_new_protected:Npn \prop_map_inline:Nn #1#2
11750 {
11751     \cs_gset_eq:cn
11752     { __prop_map_ \int_use:N \g__kernel_prg_map_int :wn } \__prop_pair:wn
11753     \int_gincr:N \g__kernel_prg_map_int
11754     \cs_gset_protected:Npn \__prop_pair:wn ##1 \s__prop ##2 {#2}
11755     #1
11756     \prg_break_point:Nn \prop_map_break:
11757     {
11758         \int_gdecr:N \g__kernel_prg_map_int
11759         \cs_gset_eq:Nc \__prop_pair:wn
11760         { __prop_map_ \int_use:N \g__kernel_prg_map_int :wn }
11761     }
11762 }
11763 \cs_generate_variant:Nn \prop_map_inline:Nn { c }

```

(End definition for `\prop_map_inline:Nn`. This function is documented on page 150.)

\prop_map_tokens:Nn The mapping is very similar to `\prop_map_function:NN`. The `\use_i:nn` removes the leading `\s__prop`. The odd construction `\use:n {#1}` allows #1 to contain any token without interfering with `\prop_map_break:`. The loop stops when the argument delimited by `__prop_pair:wn` is `\prg_break:` instead of being empty.

\prop_map_tokens:cn

__prop_map_tokens:nwwn

```

11764 \cs_new:Npn \prop_map_tokens:Nn #1#2
11765 {
11766     \exp_last_unbraced:Nno
11767     \use_i:nn { \__prop_map_tokens:nwwn {#2} } #1
11768     \prg_break: \__prop_pair:wn \s__prop { } \prg_break_point:
11769     \prg_break_point:Nn \prop_map_break: { }
11770 }
11771 \cs_new:Npn \__prop_map_tokens:nwwn #1#2 \__prop_pair:wn #3 \s__prop #4
11772 {
11773     #2
11774     \use:n {#1} {#3} {#4}
11775     \__prop_map_tokens:nwwn {#1}
11776 }
11777 \cs_generate_variant:Nn \prop_map_tokens:Nn { c }

```

(End definition for `\prop_map_tokens:Nn` and `_prop_map_tokens:nwn`. This function is documented on page 150.)

`\prop_map_break:` The break statements are based on the general `\prg_map_break:Nn`.
`\prop_map_break:n`

```
11778 \cs_new:Npn \prop_map_break:
11779   { \prg_map_break:Nn \prop_map_break: { } }
11780 \cs_new:Npn \prop_map_break:n
11781   { \prg_map_break:Nn \prop_map_break: }
```

(End definition for `\prop_map_break:` and `\prop_map_break:n`. These functions are documented on page 150.)

18.7 Viewing property lists

`\prop_show:N` Apply the general `_kernel_chk_defined:NT` and `\msg_show:nnnnnn`. Contrarily to sequences and comma lists, we use `\msg_show_item:nn` to format both the key and the value for each pair.
`\prop_show:c`
`\prop_log:N`
`\prop_log:c`

```
11782 \cs_new_protected:Npn \prop_show:N { \_prop_show:NN \msg_show:nnxxxx }
11783 \cs_generate_variant:Nn \prop_show:N { c }
11784 \cs_new_protected:Npn \prop_log:N { \_prop_show:NN \msg_log:nnxxxx }
11785 \cs_generate_variant:Nn \prop_log:N { c }
11786 \cs_new_protected:Npn \_prop_show:NN #1#2
11787   {
11788     \_kernel_chk_defined:NT #2
11789     {
11790       #1 { LaTeX/kernel } { show-prop }
11791       { \token_to_str:N #2 }
11792       { \prop_map_function:NN #2 \msg_show_item:nn }
11793       { } { }
11794     }
11795   }
```

(End definition for `\prop_show:N` and `\prop_log:N`. These functions are documented on page 151.)

```
11796 \endpackage
```

19 l3msg implementation

```
11797 \begin{package}
11798 \begin{@@=msg}
```

`\l__msg_internal_tl` A general scratch for the module.

```
11799 \tl_new:N \l__msg_internal_tl
```

(End definition for `\l__msg_internal_tl`.)

`\l__msg_name_str` Used to save module info when creating messages.

```
\l__msg_text_str
11800 \str_new:N \l__msg_name_str
11801 \str_new:N \l__msg_text_str
```

(End definition for `\l__msg_name_str` and `\l__msg_text_str`.)

19.1 Internal auxiliaries

`\s__msg_mark` Internal scan marks.

```
\s__msg_stop 11802 \scan_new:N \s__msg_mark
11803 \scan_new:N \s__msg_stop
```

(End definition for `\s__msg_mark` and `\s__msg_stop`.)

`__msg_use_none_delimit_by_s_stop:w` Functions to gobble up to a scan mark.

```
11804 \cs_new:Npn \__msg_use_none_delimit_by_s_stop:w #1 \s__msg_stop { }
```

(End definition for `__msg_use_none_delimit_by_s_stop:w`.)

19.2 Creating messages

Messages are created and used separately, so there two parts to the code here. First, a mechanism for creating message text. This is pretty simple, as there is not actually a lot to do.

`\c__msg_text_prefix_tl` Locations for the text of messages.

```
\c__msg_more_text_prefix_tl 11805 \tl_const:Nn \c__msg_text_prefix_tl { msg~text~>~ }
11806 \tl_const:Nn \c__msg_more_text_prefix_tl { msg~extra~text~>~ }
```

(End definition for `\c__msg_text_prefix_tl` and `\c__msg_more_text_prefix_tl`.)

`\msg_if_exist_p:nn` Test whether the control sequence containing the message text exists or not.

```
\msg_if_exist:nnTF 11807 \prg_new_conditional:Npnn \msg_if_exist:nn #1#2 { p , T , F , TF }
11808 {
11809     \cs_if_exist:cTF { \c__msg_text_prefix_tl #1 / #2 }
11810     { \prg_return_true: } { \prg_return_false: }
11811 }
```

(End definition for `\msg_if_exist:nnTF`. This function is documented on page 154.)

`__msg_chk_if_free:nn` This auxiliary is similar to `__kernel_chk_if_free_cs:N`, and is used when defining messages with `\msg_new:nnnn`.

```
11812 \cs_new_protected:Npn \__msg_chk_free:nn #1#2
11813 {
11814     \msg_if_exist:nnT {#1} {#2}
11815     {
11816         \__kernel_msg_error:nxxx { kernel } { message-already-defined }
11817         {#1} {#2}
11818     }
11819 }
```

(End definition for `__msg_chk_if_free:nn`.)

`\msg_new:nnnn` Setting a message simply means saving the appropriate text into two functions. A sanity check first.

```
\msg_new:nnn 11820 \cs_new_protected:Npn \msg_new:nnnn #1#2
\msg_gset:nnnn 11821 {
\msg_gset:nnn 11822     \__msg_chk_free:nn {#1} {#2}
\msg_set:nnnn 11823     \msg_gset:nnnn {#1} {#2}
\msg_set:nnn 11824 }
11825 \cs_new_protected:Npn \msg_new:nnn #1#2#3
```

```

11826 { \msg_new:nnnn {#1} {#2} {#3} { } }
11827 \cs_new_protected:Npn \msg_set:nnnn #1#2#3#4
11828 {
11829   \cs_set:cpn { \c__msg_text_prefix_tl #1 / #2 }
11830   ##1##2##3##4 {#3}
11831   \cs_set:cpn { \c__msg_more_text_prefix_tl #1 / #2 }
11832   ##1##2##3##4 {#4}
11833 }
11834 \cs_new_protected:Npn \msg_set:nnn #1#2#3
11835 { \msg_set:nnnn {#1} {#2} {#3} { } }
11836 \cs_new_protected:Npn \msg_gset:nnnn #1#2#3#4
11837 {
11838   \cs_gset:cpn { \c__msg_text_prefix_tl #1 / #2 }
11839   ##1##2##3##4 {#3}
11840   \cs_gset:cpn { \c__msg_more_text_prefix_tl #1 / #2 }
11841   ##1##2##3##4 {#4}
11842 }
11843 \cs_new_protected:Npn \msg_gset:nnn #1#2#3
11844 { \msg_gset:nnnn {#1} {#2} {#3} { } }

```

(End definition for `\msg_new:nnnn` and others. These functions are documented on page [153](#).)

19.3 Messages: support functions and text

```

\c__msg_coding_error_text_tl Simple pieces of text for messages.
\c__msg_continue_text_tl
\c__msg_critical_text_tl
\c__msg_fatal_text_tl
\c__msg_help_text_tl
\c__msg_no_info_text_tl
\c__msg_on_line_text_tl
\c__msg_return_text_tl
\c__msg_trouble_text_tl
11845 \tl_const:Nn \c__msg_coding_error_text_tl
11846 {
11847   This~is~a~coding~error.
11848   \ \ \
11849 }
11850 \tl_const:Nn \c__msg_continue_text_tl
11851 { Type~<return>~to~continue }
11852 \tl_const:Nn \c__msg_critical_text_tl
11853 { Reading~the~current~file~'\g_file_curr_name_str'~will~stop. }
11854 \tl_const:Nn \c__msg_fatal_text_tl
11855 { This~is~a~fatal~error:~LaTeX~will~abort. }
11856 \tl_const:Nn \c__msg_help_text_tl
11857 { For~immediate~help~type~H~<return> }
11858 \tl_const:Nn \c__msg_no_info_text_tl
11859 {
11860   LaTeX~does~not~know~anything~more~about~this~error,~sorry.
11861   \c__msg_return_text_tl
11862 }
11863 \tl_const:Nn \c__msg_on_line_text_tl { on~line }
11864 \tl_const:Nn \c__msg_return_text_tl
11865 {
11866   \ \ \
11867   Try~typing~<return>~to~proceed.
11868   \ \
11869   If~that~doesn't~work,~type~X~<return>~to~quit.
11870 }
11871 \tl_const:Nn \c__msg_trouble_text_tl
11872 {
11873   \ \ \

```

```

11874     More~errors~will~almost~certainly~follow: \\
11875     the~LaTeX~run~should~be~aborted.
11876 }

```

(End definition for `\c_msg_coding_error_text_tl` and others.)

`\msg_line_number:` For writing the line number nicely. `\msg_line_context:` was set up earlier, so this is not new.

```

\msg_line_context:
11877 \cs_new:Npn \msg_line_number: { \int_use:N \tex_inputlineno:D }
11878 \cs_gset:Npn \msg_line_context:
11879 {
11880     \c_msg_on_line_text_tl
11881     \c_space_tl
11882     \msg_line_number:
11883 }

```

(End definition for `\msg_line_number:` and `\msg_line_context:`. These functions are documented on page 154.)

19.4 Showing messages: low level mechanism

```

\__msg_interrupt:Nnnn
\__msg_no_more_text:nnnn

```

The low-level interruption macro is rather opaque, unfortunately. Depending on the availability of more information there is a choice of how to set up the further help. We feed the extra help text and the message itself to a wrapping auxiliary, in this order because we must first setup TeX's `\errhelp` register before issuing an `\errmessage`. To deal with the various cases of critical or fatal errors with and without help text, there is a bit of argument-passing to do.

```

11884 \cs_new_protected:Npn \__msg_interrupt:NnnnN #1#2#3#4#5
11885 {
11886     \str_set:Nx \l__msg_text_str { #1 {#2} }
11887     \str_set:Nx \l__msg_name_str { \msg_module_name:n {#2} }
11888     \cs_if_eq:cNTF
11889     { \c__msg_more_text_prefix_tl #2 / #3 }
11890     \__msg_no_more_text:nnnn
11891     {
11892         \__msg_interrupt_wrap:nnn
11893         { \use:c { \c__msg_text_prefix_tl #2 / #3 } #4 }
11894         { \c__msg_continue_text_tl }
11895         {
11896             \c__msg_no_info_text_tl
11897             \tl_if_empty:NF #5
11898             { \\ \\ #5 }
11899         }
11900     }
11901     {
11902         \__msg_interrupt_wrap:nnn
11903         { \use:c { \c__msg_text_prefix_tl #2 / #3 } #4 }
11904         { \c__msg_help_text_tl }
11905         {
11906             \use:c { \c__msg_more_text_prefix_tl #2 / #3 } #4
11907             \tl_if_empty:NF #5
11908             { \\ \\ #5 }
11909         }
11910     }
}

```

```

11911 }
11912 \cs_new:Npn \__msg_no_more_text:nnnn #1#2#3#4 { }

(End definition for \__msg_interrupt:Nnnn and \__msg_no_more_text:nnnn.)

```

```

\__msg_interrupt_wrap:nnn
\__msg_interrupt_text:n
\__msg_interrupt_more_text:n

```

First setup TeX's `\errhelp` register with the extra help #1, then build a nice-looking error message with #2. Everything is done using x-type expansion as the new line markers are different for the two type of text and need to be correctly set up. The auxiliary `__msg_interrupt_more_text:n` receives its argument as a line-wrapped string, which is thus unaffected by expansion. We have to split the main text into two parts as only the “message” itself is wrapped with a leader: the generic help is wrapped at full width. We also have to allow for the two characters used by `\errmessage` itself.

```

11913 \cs_new_protected:Npn \__msg_interrupt_wrap:nnn #1#2#3
11914 {
11915   \iow_wrap:nnnN { \ \ #3 } { } { } \__msg_interrupt_more_text:n
11916   \group_begin:
11917     \int_sub:Nn \l_iow_line_count_int { 2 }
11918     \iow_wrap:nxnN { \l__msg_text_str : ~ #1 }
11919     {
11920       ( \l__msg_name_str )
11921       \prg_replicate:nn
11922       {
11923         \str_count:N \l__msg_text_str
11924         - \str_count:N \l__msg_name_str
11925         + 2
11926       }
11927       { ~ }
11928     }
11929     { } \__msg_interrupt_text:n
11930     \iow_wrap:nnnN { \l__msg_internal_tl \ \ \ #2 } { } { }
11931     \__msg_interrupt:n
11932   }
11933 \cs_new_protected:Npn \__msg_interrupt_text:n #1
11934 {
11935   \group_end:
11936   \tl_set:Nn \l__msg_internal_tl {#1}
11937 }
11938 \cs_new_protected:Npn \__msg_interrupt_more_text:n #1
11939 { \exp_args:Nx \tex_errhelp:D { #1 \iow_newline: } }

```

(End definition for `__msg_interrupt_wrap:nnn`, `__msg_interrupt_text:n`, and `__msg_interrupt_more_text:n`.)

```
\__msg_interrupt:n
```

The business end of the process starts by producing some visual separation of the message from the main part of the log. The error message needs to be printed with everything made “invisible”: TeX's own information involves the macro in which `\errmessage` is called, and the end of the argument of the `\errmessage`, including the closing brace. We use an active `!` to call the `\errmessage` primitive, and end its argument with `\use_none:n {<spaces>}` which fills the output with spaces. Two trailing closing braces are turned into spaces to hide them as well. The group in which we alter the definition of the active `!` is closed before producing the message: this ensures that tokens inserted by typing `I` in the command-line are inserted after the message is entirely cleaned up.

The `__kernel_iow_with:Nnn` auxiliary, defined in `l3file`, expects an *integer variable*, an integer *value*, and some *code*. It runs the *code* after ensuring that the

$\langle integer\ variable \rangle$ takes the given $\langle value \rangle$, then restores the former value of the $\langle integer\ variable \rangle$ if needed. We use it to ensure that the `\newlinechar` is 10, as needed for `\iow_newline:` to work, and that `\errorcontextlines` is -1 , to avoid showing irrelevant context. Note that restoring the former value of these integers requires inserting tokens after the `\errmessage`, which go in the way of tokens which could be inserted by the user. This is unavoidable.

```

11940 \group_begin:
11941   \char_set_lccode:nn { 38 } { 32 } % &
11942   \char_set_lccode:nn { 46 } { 32 } % .
11943   \char_set_lccode:nn { 123 } { 32 } % {
11944   \char_set_lccode:nn { 125 } { 32 } % }
11945   \char_set_catcode_active:N \&
11946 \tex_lowercase:D
11947 {
11948   \group_end:
11949   \cs_new_protected:Npn \_msg_interrupt:n #1
11950   {
11951     \iow_term:n { }
11952     \_kernel_iow_with:Nnn \tex_newlinechar:D { ^^J }
11953     {
11954       \_kernel_iow_with:Nnn \tex_errorcontextlines:D { -1 }
11955       {
11956         \group_begin:
11957         \cs_set_protected:Npn &
11958         {
11959           \tex_errmessage:D
11960           {
11961             #1
11962             \use_none:n
11963             { ..... }
11964           }
11965         }
11966         \exp_after:wN
11967         \group_end:
11968         &
11969       }
11970     }
11971   }
11972 }

```

(End definition for `_msg_interrupt:n`.)

19.5 Displaying messages

L^AT_EX is handling error messages and so the T_EX ones are disabled.

```

11973 \int_gset:Nn \tex_errorcontextlines:D { -1 }

```

```

\msg_fatal_text:n
\msg_critical_text:n
\msg_error_text:n
\msg_warning_text:n
\msg_info_text:n
\_msg_text:nn
\_msg_text:n

```

A function for issuing messages: both the text and order could in principle vary. The module name may be empty for kernel messages, hence the slightly contorted code path for a space.

```

11974 \cs_new:Npn \msg_fatal_text:n #1
11975 {
11976   Fatal ~

```

```

11977     \msg_error_text:n {#1}
11978   }
11979   \cs_new:Npn \msg_critical_text:n #1
11980   {
11981     Critical ~
11982     \msg_error_text:n {#1}
11983   }
11984   \cs_new:Npn \msg_error_text:n #1
11985   { \__msg_text:nn {#1} { Error } }
11986   \cs_new:Npn \msg_warning_text:n #1
11987   { \__msg_text:nn {#1} { Warning } }
11988   \cs_new:Npn \msg_info_text:n #1
11989   { \__msg_text:nn {#1} { Info } }
11990   \cs_new:Npn \__msg_text:nn #1#2
11991   {
11992     \exp_args:Nf \__msg_text:n { \msg_module_type:n {#1} }
11993     \msg_module_name:n {#1} ~
11994     #2
11995   }
11996   \cs_new:Npn \__msg_text:n #1
11997   {
11998     \tl_if_blank:nF {#1}
11999     { #1 ~ }
12000   }

```

(End definition for `\msg_fatal_text:n` and others. These functions are documented on page 154.)

`\g_msg_module_name_prop` For storing public module information: the kernel data is set up in advance.
`\g_msg_module_type_prop`

```

12001 \prop_new:N \g_msg_module_name_prop
12002 \prop_gput:Nnn \g_msg_module_name_prop { LaTeX } { LaTeX3 }
12003 \prop_new:N \g_msg_module_type_prop
12004 \prop_gput:Nnn \g_msg_module_type_prop { LaTeX } { }

```

(End definition for `\g_msg_module_name_prop` and `\g_msg_module_type_prop`. These variables are documented on page 155.)

`\msg_module_type:n` Contextual footer information, with the potential to give modules an alternative name.

```

12005 \cs_new:Npn \msg_module_type:n #1
12006 {
12007   \prop_if_in:NnTF \g_msg_module_type_prop {#1}
12008   { \prop_item:Nn \g_msg_module_type_prop {#1} }
12009   { Package }
12010 }

```

(End definition for `\msg_module_type:n`. This function is documented on page 155.)

`\msg_module_name:n` Contextual footer information, with the potential to give modules an alternative name.
`\msg_see_documentation_text:n`

```

12011 \cs_new:Npn \msg_module_name:n #1
12012 {
12013   \prop_if_in:NnTF \g_msg_module_name_prop {#1}
12014   { \prop_item:Nn \g_msg_module_name_prop {#1} }
12015   {#1}
12016 }
12017 \cs_new:Npn \msg_see_documentation_text:n #1
12018 {

```

```

12019 See-the~ \msg_module_name:n {#1} ~
12020 documentation-for-further-information.
12021 }

```

(End definition for \msg_module_name:n and \msg_see_documentation_text:n. These functions are documented on page 155.)

_msg_class_new:nn

```

12022 \group_begin:
12023 \cs_set_protected:Npn \_msg_class_new:nn #1#2
12024 {
12025   \prop_new:c { l\_msg_redirect_ #1 _prop }
12026   \cs_new_protected:cpn { \_msg_ #1 _code:nnnnnn }
12027     ##1##2##3##4##5##6 {#2}
12028   \cs_new_protected:cpn { msg_ #1 :nnnnnn } ##1##2##3##4##5##6
12029     {
12030       \use:x
12031       {
12032         \exp_not:n { \_msg_use:nnnnnn {#1} {##1} {##2} }
12033         { \tl_to_str:n {##3} } { \tl_to_str:n {##4} }
12034         { \tl_to_str:n {##5} } { \tl_to_str:n {##6} }
12035       }
12036     }
12037   \cs_new_protected:cpx { msg_ #1 :nnnnn } ##1##2##3##4##5
12038     { \exp_not:c { msg_ #1 :nnnnnn } {##1} {##2} {##3} {##4} {##5} { } }
12039   \cs_new_protected:cpx { msg_ #1 :nnnn } ##1##2##3##4
12040     { \exp_not:c { msg_ #1 :nnnnnn } {##1} {##2} {##3} {##4} { } { } }
12041   \cs_new_protected:cpx { msg_ #1 :nnn } ##1##2##3
12042     { \exp_not:c { msg_ #1 :nnnnnn } {##1} {##2} {##3} { } { } { } }
12043   \cs_new_protected:cpx { msg_ #1 :nn } ##1##2
12044     { \exp_not:c { msg_ #1 :nnnnnn } {##1} {##2} { } { } { } { } }
12045   \cs_new_protected:cpx { msg_ #1 :nnxxxx } ##1##2##3##4##5##6
12046     {
12047       \use:x
12048       {
12049         \exp_not:N \exp_not:n
12050         { \exp_not:c { msg_ #1 :nnnnnn } {##1} {##2} }
12051         {##3} {##4} {##5} {##6}
12052       }
12053     }
12054   \cs_new_protected:cpx { msg_ #1 :nnxxx } ##1##2##3##4##5
12055     { \exp_not:c { msg_ #1 :nnxxxx } {##1} {##2} {##3} {##4} {##5} { } }
12056   \cs_new_protected:cpx { msg_ #1 :nnxx } ##1##2##3##4
12057     { \exp_not:c { msg_ #1 :nnxxxx } {##1} {##2} {##3} {##4} { } { } }
12058   \cs_new_protected:cpx { msg_ #1 :nnx } ##1##2##3
12059     { \exp_not:c { msg_ #1 :nnxxxx } {##1} {##2} {##3} { } { } { } }
12060 }

```

(End definition for _msg_class_new:nn.)

```

\msg_fatal:nnnnnn For fatal errors, after the error message TeX bails out. We force a bail out rather than
\msg_fatal:nnxxxx using \end as this means it does not matter if we are in a context where normally the
\msg_fatal:nnnnn run cannot end.
\msg_fatal:nnxxx
\msg_fatal:nnnn
\msg_fatal:nnxx
\msg_fatal:nnx
\msg_fatal:nn
\msg_fatal:nn
\_msg_fatal_exit:

```

```

12062 {
12063   \__msg_interrupt:NnnnN
12064   \msg_fatal_text:n {#1} {#2}
12065   { {#3} {#4} {#5} {#6} }
12066   \c__msg_fatal_text_tl
12067   \__msg_fatal_exit:
12068 }
12069 \cs_new_protected:Npn \__msg_fatal_exit:
12070 {
12071   \tex_batchmode:D
12072   \tex_read:D -1 to \l__msg_internal_tl
12073 }

```

(End definition for `\msg_fatal:nnnnnn` and others. These functions are documented on page 156.)

`\msg_critical:nnnnnn` Not quite so bad: just end the current file.

```

\msg_critical:nnxxxx 12074 \__msg_class_new:nn { critical }
\msg_critical:nnnnnn 12075 {
\msg_critical:nnxxx 12076   \__msg_interrupt:NnnnN
\msg_critical:nnnn 12077   \msg_critical_text:n {#1} {#2}
\msg_critical:nnxx 12078   { {#3} {#4} {#5} {#6} }
\msg_critical:nnn 12079   \c__msg_critical_text_tl
\msg_critical:nnx 12080   \tex_endinput:D
\msg_critical:nn 12081 }

```

(End definition for `\msg_critical:nnnnnn` and others. These functions are documented on page 156.)

`\msg_error:nnnnnn` For an error, the interrupt routine is called. We check if there is a “more text” by comparing that control sequence with a permanently empty text.

```

\msg_error:nnnnnn 12082 \__msg_class_new:nn { error }
\msg_error:nnxxxx 12083 {
\msg_error:nnnnn 12084   \__msg_interrupt:NnnnN
\msg_error:nnxx 12085   \msg_error_text:n {#1} {#2}
\msg_error:nnn 12086   { {#3} {#4} {#5} {#6} }
\msg_error:nnx 12087   \c_empty_tl
\msg_error:nn 12088 }

```

(End definition for `\msg_error:nnnnnn` and others. These functions are documented on page 156.)

`\msg_warning:nnnnnn` Warnings are printed to the terminal.

```

\msg_warning:nnxxxx 12089 \__msg_class_new:nn { warning }
\msg_warning:nnnnnn 12090 {
\msg_warning:nnxxx 12091   \str_set:Nx \l__msg_text_str { \msg_warning_text:n {#1} }
\msg_warning:nnnn 12092   \str_set:Nx \l__msg_name_str { \msg_module_name:n {#1} }
\msg_warning:nnxx 12093   \iow_term:n { }
\msg_warning:nnn 12094   \iow_wrap:nxnN
\msg_warning:nnx 12095   {
12096     \l__msg_text_str : ~
12097     \use:c { \c__msg_text_prefix_tl #1 / #2 } {#3} {#4} {#5} {#6}
12098   }
12099   {
12100     ( \l__msg_name_str )
12101     \prg_replicate:nn
12102     {
12103       \str_count:N \l__msg_text_str

```

```

12104         - \str_count:N \l__msg_name_str
12105     }
12106     { ~ }
12107 }
12108 { } \iow_term:n
12109 \iow_term:n { }
12110 }

```

(End definition for `\msg_warning:nnnnnn` and others. These functions are documented on page 156.)

`\msg_info:nnnnnn` Information only goes into the log.

```

\msg_info:nnxxxx 12111 \__msg_class_new:nn { info }
\msg_info:nnnnnn 12112 {
\msg_info:nnxxxx 12113   \str_set:Nx \l__msg_text_str { \msg_info_text:n {#1} }
\msg_info:nnnn 12114   \str_set:Nx \l__msg_name_str { \msg_module_name:n {#1} }
\msg_info:nnxx 12115   \iow_log:n { }
\msg_info:nnn 12116   \iow_wrap:nxnN
\msg_info:nnx 12117   {
\msg_info:nn 12118     \l__msg_text_str : ~
12119     \use:c { \c__msg_text_prefix_tl #1 / #2 } {#3} {#4} {#5} {#6}
12120   }
12121   {
12122     ( \l__msg_name_str )
12123     \prg_replicate:nn
12124       {
12125         \str_count:N \l__msg_text_str
12126         - \str_count:N \l__msg_name_str
12127       }
12128     { ~ }
12129   }
12130   { } \iow_log:n
12131   \iow_log:n { }
12132 }

```

(End definition for `\msg_info:nnnnnn` and others. These functions are documented on page 157.)

`\msg_log:nnnnnn` “Log” data is very similar to information, but with no extras added.

```

\msg_log:nnxxxx 12133 \__msg_class_new:nn { log }
\msg_log:nnnnnn 12134 {
\msg_log:nnxxxx 12135   \iow_wrap:nnnN
\msg_log:nnnn 12136   { \use:c { \c__msg_text_prefix_tl #1 / #2 } {#3} {#4} {#5} {#6} }
\msg_log:nnxx 12137   { } { } \iow_log:n
\msg_log:nnn 12138   }
\msg_log:nnx (End definition for \msg_log:nnnnnn and others. These functions are documented on page 157.)
\msg_log:nn

```

`\msg_term:nnnnnn` “Term” is used for communicating with the user through the terminal, like diagnostic messages, and debugging. This is similar to “log” messages, but uses the terminal output.

```

\msg_term:nnxxxx 12139 \__msg_class_new:nn { term }
\msg_term:nnxxxx 12140 {
\msg_term:nnnn 12141   \iow_wrap:nnnN
\msg_term:nnxx 12142   { \use:c { \c__msg_text_prefix_tl #1 / #2 } {#3} {#4} {#5} {#6} }
\msg_term:nnn 12143   { } { } \iow_term:n
\msg_term:nnx 12144   }
\msg_term:nn

```

(End definition for `\msg_term:nnnnnn` and others. These functions are documented on page 157.)

`\msg_none:nnnnnn`
`\msg_none:nnxxxx`
`\msg_none:nnnnnn`
`\msg_none:nnxxx`

The `none` message type is needed so that input can be gobbled.

```
12145 \__msg_class_new:nn { none } { }
```

(End definition for `\msg_none:nnnnnn` and others. These functions are documented on page 157.)

`\msg_none:nnnn`
`\msg_show:nnnnnn`
`\msg_none:nnxx`
`\msg_show:nnxxxx`
`\msg_none:nnn`
`\msg_show:nnnnnn`
`\msg_none:nnx`
`\msg_show:nnxxx`
`\msg_none:nn`
`\msg_show:nnnn`
`\msg_show:nnxx`
`\msg_show:nnn`
`\msg_show:nnx`
`\msg_show:nn`
`__msg_show:n`
`__msg_show:w`
`__msg_show_dot:w`
`__msg_show:nn`

The `show` message type is used for `\seq_show:N` and similar complicated data structures. Wrap the given text with a trailing dot (important later) then pass it to `__msg_show:n`. If there is `\>~` (or if the whole thing starts with `>~`) we split there, print the first part and show the second part using `\showtokens` (the `\exp_after:wN` ensure a nice display). Note that this primitive adds a leading `>~` and trailing dot. That is why we included a trailing dot before wrapping and removed it afterwards. If there is no `\>~` do the same but with an empty second part which adds a spurious but inevitable `>~`.

```
12146 \__msg_class_new:nn { show }
12147 {
12148   \iow_wrap:nnnN
12149   { \use:c { \c__msg_text_prefix_tl #1 / #2 } {#3} {#4} {#5} {#6} }
12150   { } { } \__msg_show:n
12151 }
12152 \cs_new_protected:Npn \__msg_show:n #1
12153 {
12154   \tl_if_in:nnTF { ^^J #1 } { ^^J > ~ }
12155   {
12156     \tl_if_in:nnTF { #1 \s__msg_mark } { . \s__msg_mark }
12157     { \__msg_show_dot:w } { \__msg_show:w }
12158     ^^J #1 \s__msg_stop
12159   }
12160   { \__msg_show:nn { ? #1 } { } }
12161 }
12162 \cs_new:Npn \__msg_show_dot:w #1 ^^J > ~ #2 . \s__msg_stop
12163 { \__msg_show:nn {#1} {#2} }
12164 \cs_new:Npn \__msg_show:w #1 ^^J > ~ #2 \s__msg_stop
12165 { \__msg_show:nn {#1} {#2} }
12166 \cs_new_protected:Npn \__msg_show:nn #1#2
12167 {
12168   \tl_if_empty:nF {#1}
12169   { \exp_args:No \iow_term:n { \use_none:n #1 } }
12170   \tl_set:Nn \l__msg_internal_tl {#2}
12171   \__kernel_iow_with:Nnn \tex_newlinechar:D { 10 }
12172   {
12173     \__kernel_iow_with:Nnn \tex_errorcontextlines:D { -1 }
12174     {
12175       \tex_showtokens:D \exp_after:wN \exp_after:wN \exp_after:wN
12176       { \exp_after:wN \l__msg_internal_tl }
12177     }
12178   }
12179 }
```

(End definition for `\msg_show:nnnnnn` and others. These functions are documented on page 271.)

End the group to eliminate `__msg_class_new:nn`.

```
12180 \group_end:
```

`__msg_class_chk_exist:nT` Checking that a message class exists. We build this from `\cs_if_free:cTF` rather than `\cs_if_exist:cTF` because that avoids reading the second argument earlier than necessary.

```

12181 \cs_new:Npn \__msg_class_chk_exist:nT #1
12182 {
12183     \cs_if_free:cTF { __msg_ #1 _code:nnnnnn }
12184     { \__kernel_msg_error:nnx { kernel } { message-class-unknown } {#1} }
12185 }
```

(End definition for __msg_class_chk_exist:nT.)

`\l__msg_class_tl` Support variables needed for the redirection system.
`\l__msg_current_class_tl`

```

12186 \tl_new:N \l__msg_class_tl
12187 \tl_new:N \l__msg_current_class_tl
```

(End definition for \l__msg_class_tl and \l__msg_current_class_tl.)

`\l__msg_redirect_prop` For redirection of individually-named messages

```

12188 \prop_new:N \l__msg_redirect_prop
```

(End definition for \l__msg_redirect_prop.)

`\l__msg_hierarchy_seq` During redirection, split the message name into a sequence: `{/module/submodule}`, `{/module}`, and `{}`.

```

12189 \seq_new:N \l__msg_hierarchy_seq
```

(End definition for \l__msg_hierarchy_seq.)

`\l__msg_class_loop_seq` Classes encountered when following redirections to check for loops.

```

12190 \seq_new:N \l__msg_class_loop_seq
```

(End definition for \l__msg_class_loop_seq.)

`__msg_use:nnnnnn` Actually using a message is a multi-step process. First, some safety checks on the message
`__msg_use_redirect_name:n` and class requested. The code and arguments are then stored to avoid passing them
`__msg_use_hierarchy:nwN` around. The assignment to `__msg_use_code:` is similar to `\tl_set:Nn`. The message
`__msg_use_redirect_module:n` is eventually produced with whatever `\l__msg_class_tl` is when `__msg_use_code:` is
`__msg_use_code:` called. Here is also a good place to suppress tracing output if the `trace` package is loaded
since all (non-expandable) messages go through this auxiliary.

```

12191 \cs_new_protected:Npn \__msg_use:nnnnnn #1#2#3#4#5#6#7
12192 {
12193     \cs_if_exist_use:N \conditionally@traceoff
12194     \msg_if_exist:nnTF {#2} {#3}
12195     {
12196         \__msg_class_chk_exist:nT {#1}
12197         {
12198             \tl_set:Nn \l__msg_current_class_tl {#1}
12199             \cs_set_protected:Npx \__msg_use_code:
12200             {
12201                 \exp_not:n
12202                 {
12203                     \use:c { __msg_ \l__msg_class_tl _code:nnnnnn }
12204                     {#2} {#3} {#4} {#5} {#6} {#7}
12205                 }
12206             }
12207         }
12208     }
12209 }
```

```

12206         }
12207         \__msg_use_redirect_name:n { #2 / #3 }
12208     }
12209 }
12210 { \__kernel_msg_error:nnxx { kernel } { message-unknown } {#2} {#3} }
12211 \cs_if_exist_use:N \conditionally@traceon
12212 }
12213 \cs_new_protected:Npn \__msg_use_code: { }

```

The first check is for a individual message redirection. If this applies then no further redirection is attempted. Otherwise, split the message name into $\langle module \rangle$, $\langle submodule \rangle$ and $\langle message \rangle$ (with an arbitrary number of slashes), and store $\{/module/submodule\}$, $\{/module\}$ and $\{\}$ into $\backslash l_msg_hierarchy_seq$. We then map through this sequence, applying the most specific redirection.

```

12214 \cs_new_protected:Npn \__msg_use_redirect_name:n #1
12215 {
12216     \prop_get:NnNTF \l__msg_redirect_prop { / #1 } \l__msg_class_tl
12217     { \__msg_use_code: }
12218     {
12219         \seq_clear:N \l__msg_hierarchy_seq
12220         \__msg_use_hierarchy:nwwN { }
12221         #1 \s__msg_mark \__msg_use_hierarchy:nwwN
12222         / \s__msg_mark \__msg_use_none_delimit_by_s_stop:w
12223         \s__msg_stop
12224         \__msg_use_redirect_module:n { }
12225     }
12226 }
12227 \cs_new_protected:Npn \__msg_use_hierarchy:nwwN #1#2 / #3 \s__msg_mark #4
12228 {
12229     \seq_put_left:Nn \l__msg_hierarchy_seq {#1}
12230     #4 { #1 / #2 } #3 \s__msg_mark #4
12231 }

```

At this point, the items of $\backslash l_msg_hierarchy_seq$ are the various levels at which we should look for a redirection. Redirections which are less specific than the argument of $\backslash _msg_use_redirect_module:n$ are not attempted. This argument is empty for a class redirection, $/module$ for a module redirection, *etc.* Loop through the sequence to find the most specific redirection, with module $\##1$. The loop is interrupted after testing for a redirection for $\##1$ equal to the argument $\#1$ (least specific redirection allowed). When a redirection is found, break the mapping, then if the redirection targets the same class, output the code with that class, and otherwise set the target as the new current class, and search for further redirections. Those redirections should be at least as specific as $\##1$.

```

12232 \cs_new_protected:Npn \__msg_use_redirect_module:n #1
12233 {
12234     \seq_map_inline:Nn \l__msg_hierarchy_seq
12235     {
12236         \prop_get:cnNTF { l__msg_redirect_ \l__msg_current_class_tl _prop }
12237         {##1} \l__msg_class_tl
12238         {
12239             \seq_map_break:n
12240             {
12241                 \tl_if_eq:NNTF \l__msg_current_class_tl \l__msg_class_tl
12242                 { \__msg_use_code: }

```



```

12243         {
12244             \tl_set_eq:NN \l__msg_current_class_tl \l__msg_class_tl
12245             \__msg_use_redirect_module:n {##1}
12246         }
12247     }
12248 }
12249 {
12250     \str_if_eq:nnT {##1} {#1}
12251     {
12252         \tl_set_eq:NN \l__msg_class_tl \l__msg_current_class_tl
12253         \seq_map_break:n { \__msg_use_code: }
12254     }
12255 }
12256 }
12257 }

```

(End definition for __msg_use:nnnnnnn and others.)

\msg_redirect_name:nnn Named message always use the given class even if that class is redirected further. An empty target class cancels any existing redirection for that message.

```

12258 \cs_new_protected:Npn \msg_redirect_name:nnn #1#2#3
12259 {
12260     \tl_if_empty:nTF {#3}
12261     { \prop_remove:Nn \l__msg_redirect_prop { / #1 / #2 } }
12262     {
12263         \__msg_class_chk_exist:nT {#3}
12264         { \prop_put:Nnn \l__msg_redirect_prop { / #1 / #2 } {#3} }
12265     }
12266 }

```

(End definition for \msg_redirect_name:nnn. This function is documented on page 159.)

\msg_redirect_class:nn If the target class is empty, eliminate the corresponding redirection. Otherwise, add the redirection. We must then check for a loop: as an initialization, we start by storing the initial class in \l__msg_current_class_tl.

\msg_redirect_module:nnn

__msg_redirect:nnn

__msg_redirect_loop_chk:nnn

__msg_redirect_loop_list:n

```

12267 \cs_new_protected:Npn \msg_redirect_class:nn
12268 { \__msg_redirect:nnn { } }
12269 \cs_new_protected:Npn \msg_redirect_module:nnn #1
12270 { \__msg_redirect:nnn { / #1 } }
12271 \cs_new_protected:Npn \__msg_redirect:nnn #1#2#3
12272 {
12273     \__msg_class_chk_exist:nT {#2}
12274     {
12275         \tl_if_empty:nTF {#3}
12276         { \prop_remove:cn { l__msg_redirect_ #2 _prop } {#1} }
12277         {
12278             \__msg_class_chk_exist:nT {#3}
12279             {
12280                 \prop_put:cnn { l__msg_redirect_ #2 _prop } {#1} {#3}
12281                 \tl_set:Nn \l__msg_current_class_tl {#2}
12282                 \seq_clear:N \l__msg_class_loop_seq
12283                 \__msg_redirect_loop_chk:nnn {#2} {#3} {#1}
12284             }
12285         }
12286     }
12287 }

```

```

12286     }
12287 }

```

Since multiple redirections can only happen with increasing specificity, a loop requires that all steps are of the same specificity. The new redirection can thus only create a loop with other redirections for the exact same module, #1, and not submodules. After some initialization above, follow redirections with `\l__msg_class_tl`, and keep track in `\l__msg_class_loop_seq` of the various classes encountered. A redirection from a class to itself, or the absence of redirection both mean that there is no loop. A redirection to the initial class marks a loop. To break it, we must decide which redirection to cancel. The user most likely wants the newly added redirection to hold with no further redirection. We thus remove the redirection starting from #2, target of the new redirection. Note that no message is emitted by any of the underlying functions: otherwise we may get an infinite loop because of a message from the message system itself.

```

12288 \cs_new_protected:Npn \__msg_redirect_loop_chk:nnn #1#2#3
12289 {
12290   \seq_put_right:Nn \l__msg_class_loop_seq {#1}
12291   \prop_get:cnNT { l__msg_redirect_ #1 _prop } {#3} \l__msg_class_tl
12292   {
12293     \str_if_eq:VnF \l__msg_class_tl {#1}
12294     {
12295       \tl_if_eq:NNTF \l__msg_class_tl \l__msg_current_class_tl
12296       {
12297         \prop_put:cnn { l__msg_redirect_ #2 _prop } {#3} {#2}
12298         \__kernel_msg_warning:nnxxxx
12299         { kernel } { message-redirect-loop }
12300         { \seq_item:Nn \l__msg_class_loop_seq { 1 } }
12301         { \seq_item:Nn \l__msg_class_loop_seq { 2 } }
12302         {#3}
12303         {
12304           \seq_map_function:NN \l__msg_class_loop_seq
12305             \__msg_redirect_loop_list:n
12306             { \seq_item:Nn \l__msg_class_loop_seq { 1 } }
12307         }
12308       }
12309       { \__msg_redirect_loop_chk:onn \l__msg_class_tl {#2} {#3} }
12310     }
12311   }
12312 }
12313 \cs_generate_variant:Nn \__msg_redirect_loop_chk:nnn { o }
12314 \cs_new:Npn \__msg_redirect_loop_list:n #1 { {#1} ~ => ~ }

```

(End definition for `\msg_redirect_class:nn` and others. These functions are documented on page 159.)

19.6 Kernel-specific functions

The kernel needs some messages of its own. These are created using pre-built functions. Two functions are provided: one more general and one which only has the short text part.

```

\__kernel_msg_new:nnnn
\__kernel_msg_new:nnn
\__kernel_msg_set:nnnn
\__kernel_msg_set:nnn
12315 \cs_new_protected:Npn \__kernel_msg_new:nnnn #1#2
12316 { \msg_new:nnnn { LaTeX } { #1 / #2 } }
12317 \cs_new_protected:Npn \__kernel_msg_new:nnn #1#2
12318 { \msg_new:nnn { LaTeX } { #1 / #2 } }

```

```

12319 \cs_new_protected:Npn \__kernel_msg_set:nnnn #1#2
12320 { \msg_set:nnnn { LaTeX } { #1 / #2 } }
12321 \cs_new_protected:Npn \__kernel_msg_set:nnn #1#2
12322 { \msg_set:nnn { LaTeX } { #1 / #2 } }

```

(End definition for __kernel_msg_new:nnnn and others.)

```

\__msg_kernel_class_new:nN
  \__msg_kernel_class_new_aux:nN

```

All the functions for kernel messages come in variants ranging from 0 to 4 arguments. Those with less than 4 arguments are defined in terms of the 4-argument variant, in a way very similar to __msg_class_new:nn. This auxiliary is destroyed at the end of the group.

```

12323 \group_begin:
12324   \cs_set_protected:Npn \__msg_kernel_class_new:nN #1
12325   { \__msg_kernel_class_new_aux:nN { __kernel_msg_ #1 } }
12326   \cs_set_protected:Npn \__msg_kernel_class_new_aux:nN #1#2
12327   {
12328     \cs_new_protected:cpn { #1 :nnnnnn } ##1##2##3##4##5##6
12329     {
12330       \use:x
12331       {
12332         \exp_not:n { #2 { LaTeX } { ##1 / ##2 } }
12333         { \tl_to_str:n {##3} } { \tl_to_str:n {##4} }
12334         { \tl_to_str:n {##5} } { \tl_to_str:n {##6} }
12335       }
12336     }
12337     \cs_new_protected:cpx { #1 :nnnnn } ##1##2##3##4##5
12338     { \exp_not:c { #1 :nnnnnn } {##1} {##2} {##3} {##4} {##5} { } }
12339     \cs_new_protected:cpx { #1 :nnnn } ##1##2##3##4
12340     { \exp_not:c { #1 :nnnnnn } {##1} {##2} {##3} {##4} { } { } }
12341     \cs_new_protected:cpx { #1 :nnn } ##1##2##3
12342     { \exp_not:c { #1 :nnnnnn } {##1} {##2} {##3} { } { } { } }
12343     \cs_new_protected:cpx { #1 :nn } ##1##2
12344     { \exp_not:c { #1 :nnnnnn } {##1} {##2} { } { } { } { } }
12345     \cs_new_protected:cpx { #1 :nnxxx } ##1##2##3##4##5##6
12346     {
12347       \use:x
12348       {
12349         \exp_not:N \exp_not:n
12350         { \exp_not:c { #1 :nnnnnn } {##1} {##2} }
12351         {##3} {##4} {##5} {##6}
12352       }
12353     }
12354     \cs_new_protected:cpx { #1 :nnxxx } ##1##2##3##4##5
12355     { \exp_not:c { #1 :nnxxx } {##1} {##2} {##3} {##4} {##5} { } }
12356     \cs_new_protected:cpx { #1 :nnxx } ##1##2##3##4
12357     { \exp_not:c { #1 :nnxxx } {##1} {##2} {##3} {##4} { } { } }
12358     \cs_new_protected:cpx { #1 :nnx } ##1##2##3
12359     { \exp_not:c { #1 :nnxxx } {##1} {##2} {##3} { } { } { } }
12360   }

```

(End definition for __msg_kernel_class_new:nN and __msg_kernel_class_new_aux:nN.)

```

\__kernel_msg_fatal:nnnnnn
\__kernel_msg_fatal:nnxxxx
\__kernel_msg_fatal:nnnnn
\__kernel_msg_fatal:nnxxx
\__kernel_msg_fatal:nnnn
\__kernel_msg_fatal:nnxx
\__kernel_msg_fatal:nnn
\__kernel_msg_fatal:nnx
\__kernel_msg_fatal:nn
  \__kernel_msg_critical:nnnnnn
  \__kernel_msg_critical:nnxxxx
\__kernel_msg_critical:nnnnn
\__kernel_msg_critical:nnxxx

```

Neither fatal kernel errors nor kernel errors can be redirected. We directly use the code for (non-kernel) fatal errors and errors, adding the “LaTeX” module name. Three functions are already defined by l3basics; we need to undefine them to avoid errors.

```

12361 \_msg_kernel_class_new:nN { fatal } \_msg_fatal_code:nnnnnn
12362 \_msg_kernel_class_new:nN { critical } \_msg_critical_code:nnnnnn
12363 \cs_undefine:N \_kernel_msg_error:nnxx
12364 \cs_undefine:N \_kernel_msg_error:nnx
12365 \cs_undefine:N \_kernel_msg_error:nn
12366 \_msg_kernel_class_new:nN { error } \_msg_error_code:nnnnnn

```

(End definition for _kernel_msg_fatal:nnnnnn and others.)

_kernel_msg_warning:nnnnnn Kernel messages which can be redirected simply use the machinery for normal messages, with the module name “`LATEX`”.

```

\_kernel_msg_warning:nnxxxx 12367 \_msg_kernel_class_new:nN { warning } \msg_warning:nnxxxxx
\_kernel_msg_warning:nnnnn 12368 \_msg_kernel_class_new:nN { info } \msg_info:nnxxxxx
\_kernel_msg_warning:nnxxx
\_kernel_msg_warning:nnnn
\_kernel_msg_warning:nnxx
\_kernel_msg_warning:nnn
\_kernel_msg_warning:nnx
\_kernel_msg_warning:nn
\_kernel_msg_info:nnnnnn
\_kernel_msg_info:nnxxxx
\_kernel_msg_info:nnnnn
\_kernel_msg_info:nnxxx
\_kernel_msg_info:nnnn
\_kernel_msg_info:nnn
\_kernel_msg_info:nnx
\_kernel_msg_info:nn

```

(End definition for _kernel_msg_warning:nnnnnn and others.)

End the group to eliminate _msg_kernel_class_new:nN.

```
12369 \group_end:
```

Error messages needed to actually implement the message system itself.

```

12370 \_kernel_msg_new:nnnn { kernel } { message-already-defined }
12371 { Message~'~#2'~for~module~'~#1'~already~defined. }
12372 {
12373   \c_msg_coding_error_text_tl
12374   LaTeX~was~asked~to~define~a~new~message~called~'~#2'~\
12375   by~the~module~'~#1'~:~this~message~already~exists.
12376   \c_msg_return_text_tl
12377 }
12378 \_kernel_msg_new:nnnn { kernel } { message-unknown }
12379 { Unknown~message~'~#2'~for~module~'~#1'~. }
12380 {
12381   \c_msg_coding_error_text_tl
12382   LaTeX~was~asked~to~display~a~message~called~'~#2'~\
12383   by~the~module~'~#1'~:~this~message~does~not~exist.
12384   \c_msg_return_text_tl
12385 }
12386 \_kernel_msg_new:nnnn { kernel } { message-class-unknown }
12387 { Unknown~message~class~'~#1'~. }
12388 {
12389   LaTeX~has~been~asked~to~redirect~messages~to~a~class~'~#1'~:\
12390   this~was~never~defined.
12391   \c_msg_return_text_tl
12392 }
12393 \_kernel_msg_new:nnnn { kernel } { message-redirect-loop }
12394 {
12395   Message~redirection~loop~caused~by~ {#1} ~>~ {#2}
12396   \tl_if_empty:nF {#3} { ~for~module~' \use_none:n #3 ' } .
12397 }
12398 {
12399   Adding~the~message~redirection~ {#1} ~>~ {#2}
12400   \tl_if_empty:nF {#3} { ~for~the~module~' \use_none:n #3 ' } ~
12401   created~an~infinite~loop~\
12402   \iow_indent:n { #4 \
12403 }

```

Messages for earlier kernel modules plus a few for l3keys which cover coding errors.

```

12404 \__kernel_msg_new:nnnn { kernel } { bad-number-of-arguments }
12405 { Function~'#1'~cannot~be~defined~with~#2~arguments. }
12406 {
12407   \c__msg_coding_error_text_tl
12408   LaTeX~has~been~asked~to~define~a~function~'#1'~with~
12409   #2~arguments.~
12410   TeX~allows~between~0~and~9~arguments~for~a~single~function.
12411 }
12412 \__kernel_msg_new:nnn { kernel } { char-active }
12413 { Cannot~generate~active~chars. }
12414 \__kernel_msg_new:nnn { kernel } { char-invalid-catcode }
12415 { Invalid~catcode~for~char~generation. }
12416 \__kernel_msg_new:nnn { kernel } { char-null-space }
12417 { Cannot~generate~null~char~as~a~space. }
12418 \__kernel_msg_new:nnn { kernel } { char-out-of-range }
12419 { Charcode~requested~out~of~engine~range. }
12420 \__kernel_msg_new:nnn { kernel } { char-space }
12421 { Cannot~generate~space~chars. }
12422 \__kernel_msg_new:nnnn { kernel } { command-already-defined }
12423 { Control~sequence~#1~already~defined. }
12424 {
12425   \c__msg_coding_error_text_tl
12426   LaTeX~has~been~asked~to~create~a~new~control~sequence~'#1'~
12427   but~this~name~has~already~been~used~elsewhere. \\ \\
12428   The~current~meaning~is:\\
12429   \\ #2
12430 }
12431 \__kernel_msg_new:nnnn { kernel } { command-not-defined }
12432 { Control~sequence~#1~undefined. }
12433 {
12434   \c__msg_coding_error_text_tl
12435   LaTeX~has~been~asked~to~use~a~control~sequence~'#1':\\
12436   this~has~not~been~defined~yet.
12437 }
12438 \__kernel_msg_new:nnnn { kernel } { empty-search-pattern }
12439 { Empty~search~pattern. }
12440 {
12441   \c__msg_coding_error_text_tl
12442   LaTeX~has~been~asked~to~replace~an~empty~pattern~by~'#1':~that~
12443   would~lead~to~an~infinite~loop!
12444 }
12445 \__kernel_msg_new:nnnn { kernel } { out-of-registers }
12446 { No~room~for~a~new~#1. }
12447 {
12448   TeX~only~supports~\int_use:N \c_max_register_int \ %
12449   of~each~type.~All~the~#1~registers~have~been~used.~
12450   This~run~will~be~aborted~now.
12451 }
12452 \__kernel_msg_new:nnnn { kernel } { non-base-function }
12453 { Function~'#1'~is~not~a~base~function }
12454 {
12455   \c__msg_coding_error_text_tl
12456   Functions~defined~through~\iow_char:N\\cs_new:Nn~must~have~

```

```

12457     a~signature~consisting~of~only~normal~arguments~'N'~and~'n'.~
12458     To~define~variants~use~\iow_char:N\cs_generate_variant:Nn~
12459     and~to~define~other~functions~use~\iow_char:N\cs_new:Npn.
12460 }
12461 \__kernel_msg_new:nnnn { kernel } { missing-colon }
12462 { Function~'~#1'~contains~no~':'. }
12463 {
12464     \c__msg_coding_error_text_tl
12465     Code~level~functions~must~contain~': '~to~separate~the~
12466     argument~specification~from~the~function~name.~This~is~
12467     needed~when~defining~conditionals~or~variants,~or~when~building~a~
12468     parameter~text~from~the~number~of~arguments~of~the~function.
12469 }
12470 \__kernel_msg_new:nnnn { kernel } { overflow }
12471 { Integers~larger~than~2^{30}~1~cannot~be~stored~in~arrays. }
12472 {
12473     An~attempt~was~made~to~store~#3~
12474     \tl_if_empty:nF {#2} { at~position~#2~ } in~the~array~'~#1'.~
12475     The~largest~allowed~value~#4~will~be~used~instead.
12476 }
12477 \__kernel_msg_new:nnnn { kernel } { out-of-bounds }
12478 { Access~to~an~entry~beyond~an~array's~bounds. }
12479 {
12480     An~attempt~was~made~to~access~or~store~data~at~position~#2~of~the~
12481     array~'~#1',~but~this~array~has~entries~at~positions~from~1~to~#3.
12482 }
12483 \__kernel_msg_new:nnnn { kernel } { protected-predicate }
12484 { Predicate~'~#1'~must~be~expandable. }
12485 {
12486     \c__msg_coding_error_text_tl
12487     LaTeX~has~been~asked~to~define~'~#1'~as~a~protected~predicate.~
12488     Only~expandable~tests~can~have~a~predicate~version.
12489 }
12490 \__kernel_msg_new:nnn { kernel } { randint-backward-range }
12491 { Bounds~ordered~backwards~in~\iow_char:N\int_rand:nn~{~#1~}~{~#2~}. }
12492 \__kernel_msg_new:nnnn { kernel } { conditional-form-unknown }
12493 { Conditional~form~'~#1'~for~function~'~#2'~unknown. }
12494 {
12495     \c__msg_coding_error_text_tl
12496     LaTeX~has~been~asked~to~define~the~conditional~form~'~#1'~of~
12497     the~function~'~#2',~but~only~'TF',~'T',~'F',~and~'p'~forms~exist.
12498 }
12499 \__kernel_msg_new:nnnn { kernel } { key-no-property }
12500 { No~property~given~in~definition~of~key~'~#1'. }
12501 {
12502     \c__msg_coding_error_text_tl
12503     Inside~\keys_define:nn~each~key~name~
12504     needs~a~property:~\\~\\
12505     \iow_indent:n { #1 .<property> } \\ \\
12506     LaTeX~did~not~find~a~'.'~to~indicate~the~start~of~a~property.
12507 }
12508 \__kernel_msg_new:nnnn { kernel } { key-property-boolean-values-only }
12509 { The~property~'~#1'~accepts~boolean~values~only. }
12510 {

```

```

12511 \c_msg_coding_error_text_tl
12512 The~property~'#1'~only~accepts~the~values~'true'~and~'false'.
12513 }
12514 \__kernel_msg_new:nnnn { kernel } { key-property-requires-value }
12515 { The~property~'#1'~requires~a~value. }
12516 {
12517 \c_msg_coding_error_text_tl
12518 LaTeX~was~asked~to~set~property~'#1'~for~key~'#2'.\\
12519 No~value~was~given~for~the~property,~and~one~is~required.
12520 }
12521 \__kernel_msg_new:nnnn { kernel } { key-property-unknown }
12522 { The~key~property~'#1'~is~unknown. }
12523 {
12524 \c_msg_coding_error_text_tl
12525 LaTeX~has~been~asked~to~set~the~property~'#1'~for~key~'#2':~
12526 this~property~is~not~defined.
12527 }
12528 \__kernel_msg_new:nnnn { kernel } { quote-in-shell }
12529 { Quotes~in~shell~command~'#1'. }
12530 { Shell~commands~cannot~contain~quotes~("). }
12531 \__kernel_msg_new:nnnn { kernel } { invalid-quark-function }
12532 { Quark~test~function~'#1'~is~invalid. }
12533 {
12534 \c_msg_coding_error_text_tl
12535 LaTeX~has~been~asked~to~create~quark~test~function~'#1'~
12536 \tl_if_empty:nTF {#2}
12537 { but~that~name~ }
12538 { with~signature~'#2',~but~that~signature~ }
12539 is~not~valid.
12540 }
12541 \__kernel_msg_new:nnn { kernel } { invalid-quark }
12542 { Invalid~quark~variable~'#1'. }
12543 \__kernel_msg_new:nnnn { kernel } { scanmark-already-defined }
12544 { Scan~mark~#1~already~defined. }
12545 {
12546 \c_msg_coding_error_text_tl
12547 LaTeX~has~been~asked~to~create~a~new~scan~mark~'#1'~
12548 but~this~name~has~already~been~used~for~a~scan~mark.
12549 }
12550 \__kernel_msg_new:nnnn { kernel } { shuffle-too-large }
12551 { The~sequence~#1~is~too~long~to~be~shuffled~by~TeX. }
12552 {
12553 TeX~has~ \int_eval:n { \c_max_register_int + 1 } ~
12554 toks~registers:~this~only~allows~to~shuffle~up~to~
12555 \int_use:N \c_max_register_int \ items.~
12556 The~list~will~not~be~shuffled.
12557 }
12558 \__kernel_msg_new:nnnn { kernel } { variable-not-defined }
12559 { Variable~#1~undefined. }
12560 {
12561 \c_msg_coding_error_text_tl
12562 LaTeX~has~been~asked~to~show~a~variable~#1,~but~this~has~not~
12563 been~defined~yet.
12564 }

```

```

12565 \__kernel_msg_new:nnnn { kernel } { variant-too-long }
12566 { Variant-form-~'~#1'~longer-than-base-signature-of-~'~#2'. }
12567 {
12568   \c__msg_coding_error_text_tl
12569   LaTeX-has-been-asked-to-create-a-variant-of-the-function-~'~#2'~
12570   with-a-signature-starting-with-~'~#1',~but-that-is-longer-than-
12571   the-signature-(part-after-the-colon)-of-~'~#2'.
12572 }
12573 \__kernel_msg_new:nnnn { kernel } { invalid-variant }
12574 { Variant-form-~'~#1'~invalid-for-base-form-~'~#2'. }
12575 {
12576   \c__msg_coding_error_text_tl
12577   LaTeX-has-been-asked-to-create-a-variant-of-the-function-~'~#2'~
12578   with-a-signature-starting-with-~'~#1',~but-cannot-change-an-argument-
12579   from-type-~'~#3'~to-type-~'~#4'.
12580 }
12581 \__kernel_msg_new:nnnn { kernel } { invalid-exp-args }
12582 { Invalid-variant-specifier-~'~#1'~in-~'~#2'. }
12583 {
12584   \c__msg_coding_error_text_tl
12585   LaTeX-has-been-asked-to-create-an-~\iow_char:N\exp_args:N...~
12586   function-with-signature-~'~N#2'~but-~'~#1'~is-not-a-valid-argument-
12587   specifier.
12588 }
12589 \__kernel_msg_new:nnn { kernel } { deprecated-variant }
12590 {
12591   Variant-form-~'~#1'~deprecated-for-base-form-~'~#2'.~
12592   One-should-not-change-an-argument-from-type-~'~#3'~to-type-~'~#4'
12593   \str_case:nnF {#3}
12594   {
12595     { n } { :-use-a-~'\token_if_eq_charcode:NNTF #4 c v V'~variant? }
12596     { N } { :-base-form-only-accepts-a-single-token-argument. }
12597     {#4} { :-base-form-is-already-a-variant. }
12598   } { . }
12599 }

```

Some errors are only needed in package mode if debugging is enabled by one of the options `enable-debug`, `check-declarations`, `log-functions`, or on the contrary if debugging is turned off. In format mode the error is somewhat different.

```

12600 \__kernel_msg_new:nnnn { kernel } { enable-debug }
12601 { To-use-~'~#1'~set-the-~'enable-debug'~option. }
12602 {
12603   The-function-~'~#1'~will-be-ignored-because-it-can-only-work-if-
12604   some-internal-functions-in-expl3-have-been-appropriately-
12605   defined.~This-only-happens-if-one-of-the-options-
12606   'enable-debug',~'check-declarations'~or~'log-functions'~was-
12607   given-as-an-option:~see-the-main-expl3-documentation.
12608 }

```

Some errors only appear in expandable settings, hence don't need a “more-text” argument.

```

12609 \__kernel_msg_new:nnn { kernel } { bad-exp-end-f }
12610 { Misused-~\exp_end_continue_f:w or~:nw }
12611 \__kernel_msg_new:nnn { kernel } { bad-variable }
12612 { Erroneous-variable-~#1 used! }

```



```

12613 \_kernel_msg_new:nnn { kernel } { misused-sequence }
12614 { A-sequence-was-misused. }
12615 \_kernel_msg_new:nnn { kernel } { misused-prop }
12616 { A-property~list-was-misused. }
12617 \_kernel_msg_new:nnn { kernel } { negative-replication }
12618 { Negative-argument-for-\iow_char:N\prg_replicate:nn. }
12619 \_kernel_msg_new:nnn { kernel } { prop-keyval }
12620 { Missing/extra-'#1'~in-'#1'~(in-'...keyval:Nn') }
12621 \_kernel_msg_new:nnn { kernel } { unknown-comparison }
12622 { Relation-'#1'~unknown:~use=,~<,~>,~==,~!=,~<=,~>=. }
12623 \_kernel_msg_new:nnn { kernel } { zero-step }
12624 { Zero-step-size-for-step-function~#1. }
12625 \cs_if_exist:NF \tex_expanded:D
12626 {
12627 \_kernel_msg_new:nnn { kernel } { e-type }
12628 { #1 ~ in~e-type~argument }
12629 }

```

```

12630 \__kernel_msg_new:nnn { kernel } { show-clist }
12631 {
12632     The~comma~list~ \tl_if_empty:nF {#1} { #1 ~ }
12633     \tl_if_empty:nTF {#2}
12634     { is~empty \\>~ . }
12635     { contains~the~items~(without~outer~braces): #2 . }
12636 }
12637 \__kernel_msg_new:nnn { kernel } { show-intarray }
12638 { The~integer~array~#1~contains~#2~items: \\ #3 . }
12639 \__kernel_msg_new:nnn { kernel } { show-prop }
12640 {
12641     The~property~list~#1~
12642     \tl_if_empty:nTF {#2}
12643     { is~empty \\>~ . }
12644     { contains~the~pairs~(without~outer~braces): #2 . }
12645 }
12646 \__kernel_msg_new:nnn { kernel } { show-seq }
12647 {
12648     The~sequence~#1~
12649     \tl_if_empty:nTF {#2}
12650     { is~empty \\>~ . }
12651     { contains~the~items~(without~outer~braces): #2 . }
12652 }
12653 \__kernel_msg_new:nnn { kernel } { show-streams }
12654 {
12655     \tl_if_empty:nTF {#2} { No~ } { The~following~ }
12656     \str_case:nn {#1}
12657     {
12658         { ior } { input ~ }
12659         { iow } { output ~ }
12660     }
12661     streams~are~
12662     \tl_if_empty:nTF {#2} { open } { in~use: #2 . }
12663 }

```

```

12664 \__kernel_msg_new:nnnn { sys } { backend-set }
12665 { Backend-configuration-already-set. }
12666 {
12667   Run-time-backend-selection-may-only-be-carried-out-once-during-a-run.~
12668   This-second-attempt-to-set-them-will-be-ignored.
12669 }
12670 \__kernel_msg_new:nnnn { sys } { wrong-backend }
12671 { Backend-request-inconsistent-with-engine:~using~'#2'~backend. }
12672 {
12673   You-have-requested-backend~'#1',~but-this-is-not~suitable~for~use~with~the~
12674   active-engine.~LaTeX3-will-use~the~'#2'~backend~instead.
12675 }

```

19.7 Expandable errors

`__msg_expandable_error:n` In expansion only context, we cannot use the normal means of reporting errors. Instead, we feed \TeX an undefined control sequence, `\LaTeX3 error:.` It is thus interrupted, and shows the context, which thanks to the odd-looking `\use:n` is

```

<argument> \LaTeX3 error:
                The error message.

```

In other words, \TeX is processing the argument of `\use:n`, which is `\LaTeX3 error: <error message>`. Then `__msg_expandable_error:w` cleans up. In fact, there is an extra subtlety: if the user inserts tokens for error recovery, they should be kept. Thus we also use an odd space character (with category code 7) and keep tokens until that space character, dropping everything else until `\s__msg_stop`. The `\exp_end:` prevents losing braces around the user-inserted text if any, and stops the expansion of `\exp:w`. The group is used to prevent `\LaTeX3~error:` from being globally equal to `\scan_stop:.`

```

12676 \group_begin:
12677 \cs_set_protected:Npn \__msg_tmp:w #1#2
12678 {
12679   \cs_new:Npn \__msg_expandable_error:n ##1
12680   {
12681     \exp:w
12682     \exp_after:wN \exp_after:wN
12683     \exp_after:wN \__msg_expandable_error:w
12684     \exp_after:wN \exp_after:wN
12685     \exp_after:wN \exp_end:
12686     \use:n { #1 #2 ##1 } #2
12687   }
12688   \cs_new:Npn \__msg_expandable_error:w ##1 #2 ##2 #2 {##1}
12689 }
12690 \exp_args:Ncx \__msg_tmp:w { LaTeX3-error: }
12691 { \char_generate:nn { '\ } { 7 } }
12692 \group_end:

```

(End definition for `__msg_expandable_error:n` and `__msg_expandable_error:w`.)

`__kernel_msg_expandable_error:nnnnn` The command built from the csname `\c__msg_text_prefix_tl LaTeX / #1 / #2` takes four arguments and builds the error text, which is fed to `__msg_expandable_error:n` with appropriate expansion: just as for usual messages the arguments are first turned to strings, then the message is fully expanded.

```

\__kernel_msg_expandable_error:nnfff
\__kernel_msg_expandable_error:nnnnn
\__kernel_msg_expandable_error:nnfff
\__kernel_msg_expandable_error:nnnn
\__kernel_msg_expandable_error:nnfff
\__kernel_msg_expandable_error:nnnn
\__kernel_msg_expandable_error:nnfff
\__kernel_msg_expandable_error:nnnn
\__kernel_msg_expandable_error:nnfff
\__kernel_msg_expandable_error:nnnn

```

```

12693 \exp_args_generate:n { oooo }
12694 \cs_new:Npn \__kernel_msg_expandable_error:nnnnnn #1#2#3#4#5#6
12695 {
12696   \exp_args:Ne \__msg_expandable_error:n
12697   {
12698     \exp_args:Nc \exp_args:Noooo
12699     { \c_msg_text_prefix_tl LaTeX / #1 / #2 }
12700     { \tl_to_str:n {#3} }
12701     { \tl_to_str:n {#4} }
12702     { \tl_to_str:n {#5} }
12703     { \tl_to_str:n {#6} }
12704   }
12705 }
12706 \cs_new:Npn \__kernel_msg_expandable_error:nnnnn #1#2#3#4#5
12707 {
12708   \__kernel_msg_expandable_error:nnnnnn
12709   {#1} {#2} {#3} {#4} {#5} { }
12710 }
12711 \cs_new:Npn \__kernel_msg_expandable_error:nnnn #1#2#3#4
12712 {
12713   \__kernel_msg_expandable_error:nnnnnn
12714   {#1} {#2} {#3} {#4} { } { }
12715 }
12716 \cs_new:Npn \__kernel_msg_expandable_error:nnn #1#2#3
12717 {
12718   \__kernel_msg_expandable_error:nnnnnn
12719   {#1} {#2} {#3} { } { } { }
12720 }
12721 \cs_new:Npn \__kernel_msg_expandable_error:nn #1#2
12722 {
12723   \__kernel_msg_expandable_error:nnnnnn
12724   {#1} {#2} { } { } { } { }
12725 }
12726 \cs_generate_variant:Nn \__kernel_msg_expandable_error:nnnnnn { nnffff }
12727 \cs_generate_variant:Nn \__kernel_msg_expandable_error:nnnnn { nnfff }
12728 \cs_generate_variant:Nn \__kernel_msg_expandable_error:nnnn { nnff }
12729 \cs_generate_variant:Nn \__kernel_msg_expandable_error:nnn { nnf }

```

(End definition for __kernel_msg_expandable_error:nnnnnn and others.)

```

\msg_expandable_error:nnnnnn Pass to an auxiliary the message to display and the module name
\msg_expandable_error:nnffff
\msg_expandable_error:nnnnn
\msg_expandable_error:nnfff
\msg_expandable_error:nnnn
\msg_expandable_error:nnff
\msg_expandable_error:nnn
\msg_expandable_error:nnf
\msg_expandable_error:nn
  \__msg_expandable_error_module:nn
12730 \cs_new:Npn \msg_expandable_error:nnnnnn #1#2#3#4#5#6
12731 {
12732   \exp_args:Ne \__msg_expandable_error_module:nn
12733   {
12734     \exp_args:Nc \exp_args:Noooo
12735     { \c_msg_text_prefix_tl #1 / #2 }
12736     { \tl_to_str:n {#3} }
12737     { \tl_to_str:n {#4} }
12738     { \tl_to_str:n {#5} }
12739     { \tl_to_str:n {#6} }
12740   }
12741   {#1}
12742 }

```

```

12743 \cs_new:Npn \msg_expandable_error:nnnnn #1#2#3#4#5
12744 { \msg_expandable_error:nnnnnn {#1} {#2} {#3} {#4} {#5} { } }
12745 \cs_new:Npn \msg_expandable_error:nnnn #1#2#3#4
12746 { \msg_expandable_error:nnnnnn {#1} {#2} {#3} {#4} { } { } }
12747 \cs_new:Npn \msg_expandable_error:nnn #1#2#3
12748 { \msg_expandable_error:nnnnnn {#1} {#2} {#3} { } { } { } }
12749 \cs_new:Npn \msg_expandable_error:nn #1#2
12750 { \msg_expandable_error:nnnnnn {#1} {#2} { } { } { } { } }
12751 \cs_generate_variant:Nn \msg_expandable_error:nnnnnn { nnffff }
12752 \cs_generate_variant:Nn \msg_expandable_error:nnnnnn { nnfff }
12753 \cs_generate_variant:Nn \msg_expandable_error:nnnnnn { nnff }
12754 \cs_generate_variant:Nn \msg_expandable_error:nnnnnn { nnf }
12755 \cs_new:Npn \__msg_expandable_error_module:nn #1#2
12756 {
12757   \exp_after:wN \exp_after:wN
12758   \exp_after:wN \__msg_use_none_delimit_by_s_stop:w
12759   \use:n { \::error ! ~ #2 : ~ #1 } \s__msg_stop
12760 }

```

(End definition for \msg_expandable_error:nnnnnn and others. These functions are documented on page 158.)

```

12761 \</package>

```

20 l3file implementation

The following test files are used for this code: m3file001.

```

12762 \*package>

```

20.1 Input operations

```

12763 \@@=ior>

```

20.1.1 Variables and constants

\l__ior_internal_tl Used as a short-term scratch variable.

```

12764 \tl_new:N \l__ior_internal_tl

```

(End definition for \l__ior_internal_tl.)

\c__ior_term_ior Reading from the terminal (with a prompt) is done using a positive but non-existent stream number. Unlike writing, there is no concept of reading from the log.

```

12765 \int_const:Nn \c__ior_term_ior { 16 }

```

(End definition for \c__ior_term_ior.)

\g__ior_streams_seq A list of the currently-available input streams to be used as a stack.

```

12766 \seq_new:N \g__ior_streams_seq

```

(End definition for \g__ior_streams_seq.)

\l__ior_stream_tl Used to recover the raw stream number from the stack.

```

12767 \tl_new:N \l__ior_stream_tl

```

(End definition for \l__ior_stream_tl.)

`\g__ior_streams_prop` The name of the file attached to each stream is tracked in a property list. To get the correct number of reserved streams in package mode the underlying mechanism needs to be queried. For L^AT_EX 2_ε and plain T_EX this data is stored in `\count16`: with the `etex` package loaded we need to subtract 1 as the register holds the number of the next stream to use. In ConT_EXt, we need to look at `\count38` but there is no subtraction: like the original plain T_EX/L^AT_EX 2_ε mechanism it holds the value of the *last* stream allocated.

```

12768 \prop_new:N \g__ior_streams_prop
12769 \int_step_inline:nnn
12770   { 0 }
12771   {
12772     \cs_if_exist:NTF \normalend
12773     { \tex_count:D 38 ~ }
12774     {
12775       \tex_count:D 16 ~ %
12776       \cs_if_exist:NT \loccount { - 1 }
12777     }
12778   }
12779   {
12780     \prop_gput:Nnn \g__ior_streams_prop {#1} { Reserved-by-format }
12781   }

```

(End definition for `\g__ior_streams_prop`.)

20.1.2 Stream management

`\ior_new:N` Reserving a new stream is done by defining the name as equal to using the terminal.
`\ior_new:c`

```

12782 \cs_new_protected:Npn \ior_new:N #1 { \cs_new_eq:NN #1 \c__ior_term_ior }
12783 \cs_generate_variant:Nn \ior_new:N { c }

```

(End definition for `\ior_new:N`. This function is documented on page 160.)

`\g_tmpa_ior` The usual scratch space.
`\g_tmpb_ior`

```

12784 \ior_new:N \g_tmpa_ior
12785 \ior_new:N \g_tmpb_ior

```

(End definition for `\g_tmpa_ior` and `\g_tmpb_ior`. These variables are documented on page 167.)

`\ior_open:Nn` Use the conditional version, with an error if the file is not found.
`\ior_open:cn`

```

12786 \cs_new_protected:Npn \ior_open:Nn #1#2
12787   { \ior_open:NnF #1 {#2} { \__kernel_file_missing:n {#2} } }
12788 \cs_generate_variant:Nn \ior_open:Nn { c }

```

(End definition for `\ior_open:Nn`. This function is documented on page 160.)

`\l__ior_file_name_tl` Data storage.

```

12789 \tl_new:N \l__ior_file_name_tl

```

(End definition for `\l__ior_file_name_tl`.)

`\ior_open:NnTF` An auxiliary searches for the file in the \TeX , \LaTeX 2_ϵ and \LaTeX 3 paths. Then pass the file found to the lower-level function which deals with streams. The `full_name` is empty when the file is not found.

```

12790 \prg_new_protected_conditional:Npnn \ior_open:Nn #1#2 { T , F , TF }
12791 {
12792   \file_get_full_name:nNTF {#2} \l__ior_file_name_tl
12793   {
12794     \__kernel_ior_open:No #1 \l__ior_file_name_tl
12795     \prg_return_true:
12796   }
12797   { \prg_return_false: }
12798 }
12799 \prg_generate_conditional_variant:Nnn \ior_open:Nn { c } { T , F , TF }

```

(End definition for `\ior_open:NnTF`. This function is documented on page 161.)

`__ior_new:N` Streams are reserved using `\newread` before they can be managed by `ior`. To prevent `ior` from being affected by redefinitions of `\newread` (such as done by the third-party package `morewrites`), this macro is saved here under a private name. The complicated code ensures that `__ior_new:N` is not `\outer` despite plain \TeX 's `\newread` being `\outer`. For \ConTeXt , we have to deal with the fact that `\newread` works like our own: it actually checks before altering definition.

```

12800 \exp_args:NNf \cs_new_protected:Npn \__ior_new:N
12801 { \exp_args:NNc \exp_after:wN \exp_stop_f: { newread } }
12802 \cs_if_exist:NT \normalend
12803 {
12804   \cs_new_eq:NN \__ior_new_aux:N \__ior_new:N
12805   \cs_set_protected:Npn \__ior_new:N #1
12806   {
12807     \cs_undefine:N #1
12808     \__ior_new_aux:N #1
12809   }
12810 }

```

(End definition for `__ior_new:N`.)

`__kernel_ior_open:Nn` The stream allocation itself uses the fact that there is a list of all of those available.
`__kernel_ior_open:No` Life gets more complex as it's important to keep things in sync. That is done using a
`__ior_open_stream:Nn` two-part approach: any streams that have already been taken up by `ior` but are now free are tracked, so we first try those. If that fails, ask plain \TeX or \LaTeX 2_ϵ for a new stream and use that number (after a bit of conversion).

```

12811 \cs_new_protected:Npn \__kernel_ior_open:Nn #1#2
12812 {
12813   \ior_close:N #1
12814   \seq_gpop:NNTF \g__ior_streams_seq \l__ior_stream_tl
12815   { \__ior_open_stream:Nn #1 {#2} }
12816   {
12817     \__ior_new:N #1
12818     \__kernel_tl_set:Nx \l__ior_stream_tl { \int_eval:n {#1} }
12819     \__ior_open_stream:Nn #1 {#2}
12820   }
12821 }
12822 \cs_generate_variant:Nn \__kernel_ior_open:Nn { No }

```

Here, we act defensively in case LuaTeX is in use with an extensionless file name.

```

12823 \cs_new_protected:Npx \__ior_open_stream:Nn #1#2
12824 {
12825   \tex_global:D \tex_chardef:D #1 = \exp_not:N \l__ior_stream_tl \scan_stop:
12826   \prop_gput:NVn \exp_not:N \g__ior_streams_prop #1 {#2}
12827   \tex_openin:D #1
12828   \sys_if_engine luatex:TF
12829   { {#2} }
12830   { \exp_not:N \__kernel_file_name_quote:n {#2} \scan_stop: }
12831 }

```

(End definition for `__kernel_ior_open:Nn` and `__ior_open_stream:Nn`.)

`\ior_close:N` Closing a stream means getting rid of it at the TeX level and removing from the various data structures. Unless the name passed is an invalid stream number (outside the range $[0, 15]$), it can be closed. On the other hand, it only gets added to the stack if it was not already there, to avoid duplicates building up.

`\ior_close:c`

```

12832 \cs_new_protected:Npn \ior_close:N #1
12833 {
12834   \int_compare:nT { -1 < #1 < \c__ior_term_ior }
12835   {
12836     \tex_closein:D #1
12837     \prop_gremove:NV \g__ior_streams_prop #1
12838     \seq_if_in:NVF \g__ior_streams_seq #1
12839     { \seq_gpush:NV \g__ior_streams_seq #1 }
12840     \cs_gset_eq:NN #1 \c__ior_term_ior
12841   }
12842 }
12843 \cs_generate_variant:Nn \ior_close:N { c }

```

(End definition for `\ior_close:N`. This function is documented on page 161.)

`\ior_show_list:` Show the property lists, but with some “pretty printing”. See the `l3msg` module. The first argument of the message is `ior` (as opposed to `iow`) and the second is empty if no read stream is open and non-empty (the list of streams formatted using `\msg_show_item_unbraced:nn`) otherwise. The code of the message `show-streams` takes care of translating `ior/iow` to English.

`\ior_log_list:`

`__ior_list:N`

```

12844 \cs_new_protected:Npn \ior_show_list: { \__ior_list:N \msg_show:nnxxxx }
12845 \cs_new_protected:Npn \ior_log_list: { \__ior_list:N \msg_log:nnxxxx }
12846 \cs_new_protected:Npn \__ior_list:N #1
12847 {
12848   #1 { LaTeX / kernel } { show-streams }
12849   { ior }
12850   {
12851     \prop_map_function:NN \g__ior_streams_prop
12852     \msg_show_item_unbraced:nn
12853   }
12854   { } { }
12855 }

```

(End definition for `\ior_show_list:`, `\ior_log_list:`, and `__ior_list:N`. These functions are documented on page 161.)

20.1.3 Reading input

\if_eof:w The primitive conditional

```
12856 \cs_new_eq:NN \if_eof:w \tex_ifeof:D
```

(End definition for \if_eof:w. This function is documented on page 167.)

\ior_if_eof_p:N To test if some particular input stream is exhausted the following conditional is provided.
\ior_if_eof:N \underline{TF} The primitive test can only deal with numbers in the range [0,15] so we catch outliers (they are exhausted).

```
12857 \prg_new_conditional:Npnn \ior_if_eof:N #1 { p , T , F , TF }
12858 {
12859   \cs_if_exist:NTF #1
12860   {
12861     \int_compare:nTF { -1 < #1 < \c__ior_term_ior }
12862     {
12863       \if_eof:w #1
12864       \prg_return_true:
12865       \else:
12866       \prg_return_false:
12867       \fi:
12868     }
12869     { \prg_return_true: }
12870   }
12871   { \prg_return_true: }
12872 }
```

(End definition for \ior_if_eof:N \underline{TF} . This function is documented on page 164.)

\ior_get:NN And here we read from files.

```
\__ior_get:NN
\ior_get:NNTF
12873 \cs_new_protected:Npn \ior_get:NN #1#2
12874 { \ior_get:NNTF #1 #2 { \tl_set:Nn #2 { \q_no_value } } }
12875 \cs_new_protected:Npn \__ior_get:NN #1#2
12876 { \tex_read:D #1 to #2 }
12877 \prg_new_protected_conditional:Npnn \ior_get:NN #1#2 { T , F , TF }
12878 {
12879   \ior_if_eof:NNTF #1
12880   { \prg_return_false: }
12881   {
12882     \__ior_get:NN #1 #2
12883     \prg_return_true:
12884   }
12885 }
```

(End definition for \ior_get:NN, __ior_get:NN, and \ior_get:NNTF. These functions are documented on page 162.)

\ior_str_get:NN Reading as strings is a more complicated wrapper, as we wish to remove the endline character and restore it afterwards.
__ior_str_get:NN
\ior_str_get:NNTF

```
12886 \cs_new_protected:Npn \ior_str_get:NN #1#2
12887 { \ior_str_get:NNTF #1 #2 { \tl_set:Nn #2 { \q_no_value } } }
12888 \cs_new_protected:Npn \__ior_str_get:NN #1#2
12889 {
12890   \exp_args:Nno \use:n
12891   {
```



```

12892         \int_set:Nn \tex_endlinechar:D { -1 }
12893         \tex_readline:D #1 to #2
12894         \int_set:Nn \tex_endlinechar:D
12895     } { \int_use:N \tex_endlinechar:D }
12896 }
12897 \prg_new_protected_conditional:Npnn \ior_str_get:NN #1#2 { T , F , TF }
12898 {
12899     \ior_if_eof:NTF #1
12900     { \prg_return_false: }
12901     {
12902         \__ior_str_get:NN #1 #2
12903         \prg_return_true:
12904     }
12905 }

```

(End definition for `\ior_str_get:NN`, `__ior_str_get:NN`, and `\ior_str_get:NNTF`. These functions are documented on page 162.)

`\c__ior_term_noprompt_ior` For reading without a prompt.

```

12906 \int_const:Nn \c__ior_term_noprompt_ior { -1 }

```

(End definition for `\c__ior_term_noprompt_ior`.)

`\ior_get_term:nN` Getting from the terminal is better with pretty-printing.

```

\ior_str_get_term:nN
\__ior_get_term:NnN
12907 \cs_new_protected:Npn \ior_get_term:nN #1#2
12908 { \__ior_get_term:NnN \__ior_get:NN {#1} #2 }
12909 \cs_new_protected:Npn \ior_str_get_term:nN #1#2
12910 { \__ior_get_term:NnN \__ior_str_get:NN {#1} #2 }
12911 \cs_new_protected:Npn \__ior_get_term:NnN #1#2#3
12912 {
12913     \group_begin:
12914     \tex_escapechar:D = -1 \scan_stop:
12915     \tl_if_blank:NTF {#2}
12916     { \exp_args:NNc #1 \c__ior_term_noprompt_ior }
12917     { \exp_args:NNc #1 \c__ior_term_ior }
12918     {#2}
12919     \exp_args:NNNv \group_end:
12920     \tl_set:Nn #3 {#2}
12921 }

```

(End definition for `\ior_get_term:nN`, `\ior_str_get_term:nN`, and `__ior_get_term:NnN`. These functions are documented on page 270.)

`\ior_map_break:` Usual map breaking functions.

```

\ior_map_break:n
12922 \cs_new:Npn \ior_map_break:
12923 { \prg_map_break:Nn \ior_map_break: { } }
12924 \cs_new:Npn \ior_map_break:n
12925 { \prg_map_break:Nn \ior_map_break: }

```

(End definition for `\ior_map_break:` and `\ior_map_break:n`. These functions are documented on page 163.)

```

\ior_map_inline:Nn
\ior_str_map_inline:Nn
  \__ior_map_inline:NNn
  \__ior_map_inline:NNNn
\__ior_map_inline_loop:NNN

```

Mapping to an input stream can be done on either a token or a string basis, hence the set up. Within that, there is a check to avoid reading past the end of a file, hence the two applications of `\ior_if_eof:N` and its lower-level analogue `\if_eof:w`. This mapping cannot be nested with twice the same stream, as the stream has only one “current line”.

```

12926 \cs_new_protected:Npn \ior_map_inline:Nn
12927   { \__ior_map_inline:NNn \__ior_get:NN }
12928 \cs_new_protected:Npn \ior_str_map_inline:Nn
12929   { \__ior_map_inline:NNn \__ior_str_get:NN }
12930 \cs_new_protected:Npn \__ior_map_inline:NNn
12931   {
12932     \int_gincr:N \g__kernel_prg_map_int
12933     \exp_args:Nc \__ior_map_inline:NNNn
12934       { \__ior_map_ \int_use:N \g__kernel_prg_map_int :n }
12935   }
12936 \cs_new_protected:Npn \__ior_map_inline:NNNn #1#2#3#4
12937   {
12938     \cs_gset_protected:Npn #1 ##1 {#4}
12939     \ior_if_eof:NF #3 { \__ior_map_inline_loop:NNN #1#2#3 }
12940     \prg_break_point:Nn \ior_map_break:
12941       { \int_gdecr:N \g__kernel_prg_map_int }
12942   }
12943 \cs_new_protected:Npn \__ior_map_inline_loop:NNN #1#2#3
12944   {
12945     #2 #3 \l__ior_internal_tl
12946     \if_eof:w #3
12947       \exp_after:wN \ior_map_break:
12948     \fi:
12949     \exp_args:No #1 \l__ior_internal_tl
12950     \__ior_map_inline_loop:NNN #1#2#3
12951   }

```

(End definition for `\ior_map_inline:Nn` and others. These functions are documented on page 163.)

```

\ior_map_variable:NNn
\ior_str_map_variable:NNn
  \__ior_map_variable:NNNn
  \__ior_map_variable_loop:NNNn

```

Since the `TEX` primitive (`\read` or `\readline`) assigns the tokens read in the same way as a token list assignment, we simply call the appropriate primitive. The end-of-loop is checked using the primitive conditional for speed.

```

12952 \cs_new_protected:Npn \ior_map_variable:NNn
12953   { \__ior_map_variable:NNNn \ior_get:NN }
12954 \cs_new_protected:Npn \ior_str_map_variable:NNn
12955   { \__ior_map_variable:NNNn \ior_str_get:NN }
12956 \cs_new_protected:Npn \__ior_map_variable:NNNn #1#2#3#4
12957   {
12958     \ior_if_eof:NF #2 { \__ior_map_variable_loop:NNNn #1#2#3 {#4} }
12959     \prg_break_point:Nn \ior_map_break: { }
12960   }
12961 \cs_new_protected:Npn \__ior_map_variable_loop:NNNn #1#2#3#4
12962   {
12963     #1 #2 #3
12964     \if_eof:w #2
12965       \exp_after:wN \ior_map_break:
12966     \fi:
12967     #4
12968     \__ior_map_variable_loop:NNNn #1#2#3 {#4}
12969   }

```

(End definition for `\ior_map_variable:Nn` and others. These functions are documented on page 163.)

20.2 Output operations

12970 `<@@=iow>`

There is a lot of similarity here to the input operations, at least for many of the basics. Thus quite a bit is copied from the earlier material with minor alterations.

20.2.1 Variables and constants

`\c_log_iow` Here we allocate two output streams for writing to the transcript file only (`\c_log_iow`)
`\c_term_iow` and to both the terminal and transcript file (`\c_term_iow`). Recent LuaTeX provide 128 write streams; we also use `\c_term_iow` as the first non-allowed write stream so its value depends on the engine.

```
12971 \int_const:Nn \c_log_iow { -1 }
12972 \int_const:Nn \c_term_iow
12973 {
12974   \bool_lazy_and:nnTF
12975     { \sys_if_engine luatex_p: }
12976     { \int_compare_p:nNn \tex_luatexversion:D > { 80 } }
12977     { 128 }
12978     { 16 }
12979 }
```

(End definition for `\c_log_iow` and `\c_term_iow`. These variables are documented on page 167.)

`\g__iow_streams_seq` A list of the currently-available output streams to be used as a stack.

```
12980 \seq_new:N \g__iow_streams_seq
```

(End definition for `\g__iow_streams_seq`.)

`\l__iow_stream_tl` Used to recover the raw stream number from the stack.

```
12981 \tl_new:N \l__iow_stream_tl
```

(End definition for `\l__iow_stream_tl`.)

`\g__iow_streams_prop` As for reads with the appropriate adjustment of the register numbers to check on.

```
12982 \prop_new:N \g__iow_streams_prop
12983 \int_step_inline:nnn
12984 { 0 }
12985 {
12986   \cs_if_exist:NTF \normalend
12987     { \tex_count:D 39 ~ }
12988     {
12989       \tex_count:D 17 ~
12990       \cs_if_exist:NT \loccount { - 1 }
12991     }
12992 }
12993 {
12994   \prop_gput:Nnn \g__iow_streams_prop {#1} { Reserved-by~format }
12995 }
```

(End definition for `\g__iow_streams_prop`.)

20.2.2 Internal auxiliaries

`\s__iow_mark` Internal scan marks.
`\s__iow_stop` 12996 `\scan_new:N \s__iow_mark`
12997 `\scan_new:N \s__iow_stop`
(End definition for \s__iow_mark and \s__iow_stop.)

`__iow_use_i_delimit_by_s_stop:nw` Functions to gobble up to a scan mark.
12998 `\cs_new:Npn __iow_use_i_delimit_by_s_stop:nw #1 #2 \s__iow_stop {#1}`
(End definition for __iow_use_i_delimit_by_s_stop:nw.)

`\q__iow_nil` Internal quarks.
12999 `\quark_new:N \q__iow_nil`
(End definition for \q__iow_nil.)

20.3 Stream management

`\iow_new:N` Reserving a new stream is done by defining the name as equal to writing to the terminal:
`\iow_new:c` odd but at least consistent.
13000 `\cs_new_protected:Npn \iow_new:N #1 { \cs_new_eq:NN #1 \c_term_iow }`
13001 `\cs_generate_variant:Nn \iow_new:N { c }`
(End definition for \iow_new:N. This function is documented on page 160.)

`\g_tmpa_iow` The usual scratch space.
`\g_tmpb_iow` 13002 `\iow_new:N \g_tmpa_iow`
13003 `\iow_new:N \g_tmpb_iow`
(End definition for \g_tmpa_iow and \g_tmpb_iow. These variables are documented on page 167.)

`__iow_new:N` As for read streams, copy `\newwrite`, making sure that it is not `\outer`.
13004 `\exp_args:NNf \cs_new_protected:Npn __iow_new:N`
13005 `{ \exp_args:NNc \exp_after:wN \exp_stop_f: { newwrite } }`
(End definition for __iow_new:N.)

`\l__iow_file_name_tl` Data storage.
13006 `\tl_new:N \l__iow_file_name_tl`
(End definition for \l__iow_file_name_tl.)

`\iow_open:Nn` The same idea as for reading, but without the path and without the need to allow for a
`\iow_open:cn` conditional version.
`__iow_open_stream:Nn` 13007 `\cs_new_protected:Npn \iow_open:Nn #1#2`
`__iow_open_stream:NV` 13008 `{`
13009 `__kernel_tl_set:Nx \l__iow_file_name_tl`
13010 `{ __kernel_file_name_sanitiz:n {#2} }`
13011 `\iow_close:N #1`
13012 `\seq_gpop:NNTF \g__iow_streams_seq \l__iow_stream_tl`
13013 `{ __iow_open_stream:NV #1 \l__iow_file_name_tl }`
13014 `{`
13015 `__iow_new:N #1`

```

13016         \_kernel_tl_set:Nx \l__iow_stream_tl { \int_eval:n {#1} }
13017         \_iow_open_stream:NV #1 \l__iow_file_name_tl
13018     }
13019 }
13020 \cs_generate_variant:Nn \iow_open:Nn { c }
13021 \cs_new_protected:Npn \_iow_open_stream:Nn #1#2
13022 {
13023     \tex_global:D \tex_chardef:D #1 = \l__iow_stream_tl \scan_stop:
13024     \prop_gput:NVn \g__iow_streams_prop #1 {#2}
13025     \tex_immediate:D \tex_openout:D
13026         #1 \_kernel_file_name_quote:n {#2} \scan_stop:
13027 }
13028 \cs_generate_variant:Nn \_iow_open_stream:Nn { NV }

```

(End definition for `\iow_open:Nn` and `_iow_open_stream:Nn`. This function is documented on page 161.)

`\iow_close:N` Closing a stream is not quite the reverse of opening one. First, the close operation is easier than the open one, and second as the stream is actually a number we can use it directly to show that the slot has been freed up.

`\iow_close:c`

```

13029 \cs_new_protected:Npn \iow_close:N #1
13030 {
13031     \int_compare:nT { - \c_log_iow < #1 < \c_term_iow }
13032     {
13033         \tex_immediate:D \tex_closeout:D #1
13034         \prop_gremove:NV \g__iow_streams_prop #1
13035         \seq_if_in:NVF \g__iow_streams_seq #1
13036             { \seq_gpush:NV \g__iow_streams_seq #1 }
13037         \cs_gset_eq:NN #1 \c_term_iow
13038     }
13039 }
13040 \cs_generate_variant:Nn \iow_close:N { c }

```

(End definition for `\iow_close:N`. This function is documented on page 161.)

`\iow_show_list:` Done as for input, but with a copy of the auxiliary so the name is correct.

`\iow_log_list:`

`_iow_list:N`

```

13041 \cs_new_protected:Npn \iow_show_list: { \_iow_list:N \msg_show:nnxxxx }
13042 \cs_new_protected:Npn \iow_log_list: { \_iow_list:N \msg_log:nnxxxx }
13043 \cs_new_protected:Npn \_iow_list:N #1
13044 {
13045     #1 { LaTeX / kernel } { show-streams }
13046     { iow }
13047     {
13048         \prop_map_function:NN \g__iow_streams_prop
13049             \msg_show_item_unbraced:nn
13050     }
13051     { } { }
13052 }

```

(End definition for `\iow_show_list:`, `\iow_log_list:`, and `_iow_list:N`. These functions are documented on page 161.)

20.3.1 Deferred writing

`\iow_shipout_x:Nn` First the easy part, this is the primitive, which expects its argument to be braced.

```

\iow_shipout_x:Nx 13053 \cs_new_protected:Npn \iow_shipout_x:Nn #1#2
\iow_shipout_x:cn 13054 { \tex_write:D #1 {#2} }
\iow_shipout_x:cx 13055 \cs_generate_variant:Nn \iow_shipout_x:Nn { c, Nx, cx }
```

(End definition for `\iow_shipout_x:Nn`. This function is documented on page 165.)

`\iow_shipout:Nn` With ε -TeX available deferred writing without expansion is easy.

```

\iow_shipout:Nx 13056 \cs_new_protected:Npn \iow_shipout:Nn #1#2
\iow_shipout:cn 13057 { \tex_write:D #1 { \exp_not:n {#2} } }
\iow_shipout:cx 13058 \cs_generate_variant:Nn \iow_shipout:Nn { c, Nx, cx }
```

(End definition for `\iow_shipout:Nn`. This function is documented on page 165.)

20.3.2 Immediate writing

`__kernel_iow_with:Nnn` If the integer #1 is equal to #2, just leave #3 in the input stream. Otherwise, pass the old value to an auxiliary, which sets the integer to the new value, runs the code, and restores the integer.

```

13059 \cs_new_protected:Npn __kernel_iow_with:Nnn #1#2
13060 {
13061   \int_compare:nNnTF {#1} = {#2}
13062   { \use:n }
13063   { \exp_args:No __iow_with:nNnn { \int_use:N #1 } #1 {#2} }
13064 }
13065 \cs_new_protected:Npn __iow_with:nNnn #1#2#3#4
13066 {
13067   \int_set:Nn #2 {#3}
13068   #4
13069   \int_set:Nn #2 {#1}
13070 }
```

(End definition for `__kernel_iow_with:Nnn` and `__iow_with:nNnn`.)

`\iow_now:Nn` This routine writes the second argument onto the output stream without expansion. If this stream isn't open, the output goes to the terminal instead. If the first argument is no output stream at all, we get an internal error. We don't use the expansion done by `\write` to get the Nx variant, because it differs in subtle ways from x-expansion, namely, macro parameter characters would not need to be doubled. We set the `\newlinechar` to 10 using `__kernel_iow_with:Nnn` to support formats such as plain TeX: otherwise, `\iow_newline:` would not work. We do not do this for `\iow_shipout:Nn` or `\iow_shipout_x:Nn`, as TeX looks at the value of the `\newlinechar` at shipout time in those cases.

```

13071 \cs_new_protected:Npn \iow_now:Nn #1#2
13072 {
13073   __kernel_iow_with:Nnn \tex_newlinechar:D { '\^^J }
13074   { \tex_immediate:D \tex_write:D #1 { \exp_not:n {#2} } }
13075 }
13076 \cs_generate_variant:Nn \iow_now:Nn { c, Nx, cx }
```

(End definition for `\iow_now:Nn`. This function is documented on page 164.)

\iow_log:n Writing to the log and the terminal directly are relatively easy.

\iow_log:x 13077 \cs_set_protected:Npn \iow_log:x { \iow_now:Nx \c_log_iow }

\iow_term:n 13078 \cs_new_protected:Npn \iow_log:n { \iow_now:Nn \c_log_iow }

\iow_term:x 13079 \cs_set_protected:Npn \iow_term:x { \iow_now:Nx \c_term_iow }

13080 \cs_new_protected:Npn \iow_term:n { \iow_now:Nn \c_term_iow }

(End definition for \iow_log:n and \iow_term:n. These functions are documented on page 164.)

20.3.3 Special characters for writing

\iow_newline: Global variable holding the character that forces a new line when something is written to an output stream.

13081 \cs_new:Npn \iow_newline: { ^^J }

(End definition for \iow_newline:. This function is documented on page 165.)

\iow_char:N Function to write any escaped char to an output stream.

13082 \cs_new_eq:NN \iow_char:N \cs_to_str:N

(End definition for \iow_char:N. This function is documented on page 165.)

20.3.4 Hard-wrapping lines to a character count

The code here implements a generic hard-wrapping function. This is used by the messaging system, but is designed such that it is available for other uses.

\l_iow_line_count_int This is the “raw” number of characters in a line which can be written to the terminal. The standard value is the line length typically used by T_EX Live and MiK_TE_X.

13083 \int_new:N \l_iow_line_count_int

13084 \int_set:Nn \l_iow_line_count_int { 78 }

(End definition for \l_iow_line_count_int. This variable is documented on page 166.)

\l__iow_newline_tl The token list inserted to produce a new line, with the *⟨run-on text⟩*.

13085 \tl_new:N \l__iow_newline_tl

(End definition for \l__iow_newline_tl.)

\l__iow_line_target_int This stores the target line count: the full number of characters in a line, minus any part for a leader at the start of each line.

13086 \int_new:N \l__iow_line_target_int

(End definition for \l__iow_line_target_int.)

__iow_set_indent:n The **one_indent** variables hold one indentation marker and its length. The **__iow_unindent:w** auxiliary removes one indentation. The function **__iow_set_indent:n** (that could possibly be public) sets the indentation in a consistent way. We set it to four spaces by default.

13087 \tl_new:N \l__iow_one_indent_tl

13088 \int_new:N \l__iow_one_indent_int

13089 \cs_new:Npn __iow_unindent:w { }

13090 \cs_new_protected:Npn __iow_set_indent:n #1

13091 {

13092 __kernel_tl_set:Nx \l__iow_one_indent_tl

```

13093     { \exp_args:No \__kernel_str_to_other_fast:n { \tl_to_str:n {#1} } }
13094 \int_set:Nn \l__iow_one_indent_int
13095     { \str_count:N \l__iow_one_indent_tl }
13096 \exp_last_unbraced:NNo
13097     \cs_set:Npn \__iow_unindent:w \l__iow_one_indent_tl { }
13098 }
13099 \exp_args:Nx \__iow_set_indent:n { \prg_replicate:nn { 4 } { ~ } }

```

(End definition for `__iow_set_indent:n` and others.)

`\l__iow_indent_tl` The current indentation (some copies of `\l__iow_one_indent_tl`) and its number of
`\l__iow_indent_int` characters.

```

13100 \tl_new:N \l__iow_indent_tl
13101 \int_new:N \l__iow_indent_int

```

(End definition for `\l__iow_indent_tl` and `\l__iow_indent_int`.)

`\l__iow_line_tl` These hold the current line of text and a partial line to be added to it, respectively.
`\l__iow_line_part_tl`

```

13102 \tl_new:N \l__iow_line_tl
13103 \tl_new:N \l__iow_line_part_tl

```

(End definition for `\l__iow_line_tl` and `\l__iow_line_part_tl`.)

`\l__iow_line_break_bool` Indicates whether the line was broken precisely at a chunk boundary.

```

13104 \bool_new:N \l__iow_line_break_bool

```

(End definition for `\l__iow_line_break_bool`.)

`\l__iow_wrap_tl` Used for the expansion step before detokenizing, and for the output from wrapping text:
 fully expanded and with lines which are not overly long.

```

13105 \tl_new:N \l__iow_wrap_tl

```

(End definition for `\l__iow_wrap_tl`.)

`\c__iow_wrap_marker_tl` Every special action of the wrapping code is starts with the same recognizable string,
`\c__iow_wrap_end_marker_tl` `\c__iow_wrap_marker_tl`. Upon seeing that “word”, the wrapping code reads one space-
`\c__iow_wrap_newline_marker_tl` delimited argument to know what operation to perform. The setting of `\escapechar` here
`\c__iow_wrap_allow_break_marker_tl` is not very important, but makes `\c__iow_wrap_marker_tl` look marginally nicer.
`\c__iow_wrap_indent_marker_tl`
`\c__iow_wrap_unindent_marker_tl`

```

13106 \group_begin:
13107   \int_set:Nn \tex_escapechar:D { -1 }
13108   \tl_const:Nx \c__iow_wrap_marker_tl
13109     { \tl_to_str:n { \^^I \^^O \^^W \^^_ \^^W \^^R \^^A \^^P } }
13110 \group_end:
13111 \tl_map_inline:nn
13112   { { end } { newline } { allow_break } { indent } { unindent } }
13113   {
13114     \tl_const:cx { c__iow_wrap_ #1 _marker_tl }
13115     {
13116       \c__iow_wrap_marker_tl
13117       #1
13118       \c_catcode_other_space_tl
13119     }
13120   }

```

(End definition for `\c__iow_wrap_marker_tl` and others.)

`\iow_allow_break:` We set `\iow_allow_break:n` to produce an error when outside messages. Within wrapped message, it is set to `__iow_allow_break:` when valid and otherwise to `__iow_allow_break_error:`. The second produces an error expandably.

```

13121 \cs_new_protected:Npn \iow_allow_break:
13122 {
13123   \__kernel_msg_error:nnnn { kernel } { iow-indent }
13124   { \iow_wrap:nnnN } { \iow_allow_break: }
13125 }
13126 \cs_new:Npx \__iow_allow_break: { \c__iow_wrap_allow_break_marker_tl }
13127 \cs_new:Npn \__iow_allow_break_error:
13128 {
13129   \__kernel_msg_expandable_error:nnnn { kernel } { iow-indent }
13130   { \iow_wrap:nnnN } { \iow_allow_break: }
13131 }

```

(End definition for `\iow_allow_break:`, `__iow_allow_break:`, and `__iow_allow_break_error:`. This function is documented on page 269.)

`\iow_indent:n` We set `\iow_indent:n` to produce an error when outside messages. Within wrapped message, it is set to `__iow_indent:n` when valid and otherwise to `__iow_indent_error:n`. The first places the instruction for increasing the indentation before its argument, and the instruction for unindenting afterwards. The second produces an error expandably. Note that there are no forced line-break, so the indentation only changes when the next line is started.

```

13132 \cs_new_protected:Npn \iow_indent:n #1
13133 {
13134   \__kernel_msg_error:nnnnn { kernel } { iow-indent }
13135   { \iow_wrap:nnnN } { \iow_indent:n } {#1}
13136   #1
13137 }
13138 \cs_new:Npx \__iow_indent:n #1
13139 {
13140   \c__iow_wrap_indent_marker_tl
13141   #1
13142   \c__iow_wrap_unindent_marker_tl
13143 }
13144 \cs_new:Npn \__iow_indent_error:n #1
13145 {
13146   \__kernel_msg_expandable_error:nnnnn { kernel } { iow-indent }
13147   { \iow_wrap:nnnN } { \iow_indent:n } {#1}
13148   #1
13149 }

```

(End definition for `\iow_indent:n`, `__iow_indent:n`, and `__iow_indent_error:n`. This function is documented on page 166.)

`\iow_wrap:nnnN` The main wrapping function works as follows. First give `\`, `_` and other formatting commands the correct definition for messages and perform the given setup #3. The definition of `_` uses an “other” space rather than a normal space, because the latter might be absorbed by \TeX to end a number or other f-type expansions. Use `\conditionally@traceoff` if defined; it is introduced by the trace package and suppresses uninteresting tracing of the wrapping code.

```

13150 \cs_new_protected:Npn \iow_wrap:nnnN #1#2#3#4

```

```

13151 {
13152   \group_begin:
13153   \cs_if_exist_use:N \conditionally@traceoff
13154   \int_set:Nn \tex_escapechar:D { -1 }
13155   \cs_set:Npx \{ { \token_to_str:N \{ }
13156   \cs_set:Npx \# { \token_to_str:N \# }
13157   \cs_set:Npx \} { \token_to_str:N \} }
13158   \cs_set:Npx \% { \token_to_str:N \% }
13159   \cs_set:Npx \~ { \token_to_str:N \~ }
13160   \int_set:Nn \tex_escapechar:D { 92 }
13161   \cs_set_eq:NN \ \ \ \iow_newline:
13162   \cs_set_eq:NN \ \ \c_catcode_other_space_tl
13163   \cs_set_eq:NN \iow_allow_break: \__iow_allow_break:
13164   \cs_set_eq:NN \iow_indent:n \__iow_indent:n
13165   #3

```

Then fully-expand the input: in package mode, the expansion uses L^AT_EX 2_ε's `\protect` mechanism in the same way as `\typeout`. In generic mode this setting is useless but harmless. As soon as the expansion is done, reset `\iow_indent:n` to its error definition: it only works in the first argument of `\iow_wrap:nnnN`.

```

13166   \cs_set_eq:NN \protect \token_to_str:N
13167   \__kernel_tl_set:Nx \l__iow_wrap_tl {#1}
13168   \cs_set_eq:NN \iow_allow_break: \__iow_allow_break_error:
13169   \cs_set_eq:NN \iow_indent:n \__iow_indent_error:n

```

Afterwards, set the newline marker (two assignments to fully expand, then convert to a string) and initialize the target count for lines (the first line has target count `\l_iow_line_count_int` instead).

```

13170   \__kernel_tl_set:Nx \l__iow_newline_tl { \iow_newline: #2 }
13171   \__kernel_tl_set:Nx \l__iow_newline_tl { \tl_to_str:N \l__iow_newline_tl }
13172   \int_set:Nn \l__iow_line_target_int
13173   { \l_iow_line_count_int - \str_count:N \l__iow_newline_tl + 1 }

```

Sanity check.

```

13174   \int_compare:nNnT { \l__iow_line_target_int } < 0
13175   {
13176     \tl_set:Nn \l__iow_newline_tl { \iow_newline: }
13177     \int_set:Nn \l__iow_line_target_int
13178     { \l_iow_line_count_int + 1 }
13179   }

```

There is then a loop over the input, which stores the wrapped result in `\l__iow_wrap_tl`. After the loop, the resulting text is passed on to the function which has been given as a post-processor. The `\tl_to_str:N` step converts the “other” spaces back to normal spaces. The f-expansion removes a leading space from `\l__iow_wrap_tl`.

```

13180   \__iow_wrap_do:
13181   \exp_args:NNf \group_end:
13182   #4 { \tl_to_str:N \l__iow_wrap_tl }
13183   }
13184   \cs_generate_variant:Nn \iow_wrap:nnnN { nx }

```

(End definition for `\iow_wrap:nnnN`. This function is documented on page 166.)

```

\__iow_wrap_do: Escape spaces and change newlines to \c__iow_wrap_newline_marker_tl. Set up a
\__iow_wrap_fix_newline:w few variables, in particular the initial value of \l__iow_wrap_tl: the space stops the
\__iow_wrap_start:w

```

f-expansion of the main wrapping function and `\use_none:n` removes a newline marker inserted by later code. The main loop consists of repeatedly calling the `chunk` auxiliary to wrap chunks delimited by (newline or indentation) markers.

```

13185 \cs_new_protected:Npn \__iow_wrap_do:
13186 {
13187   \__kernel_tl_set:Nx \l__iow_wrap_tl
13188   {
13189     \exp_args:No \__kernel_str_to_other_fast:n \l__iow_wrap_tl
13190     \c__iow_wrap_end_marker_tl
13191   }
13192   \__kernel_tl_set:Nx \l__iow_wrap_tl
13193   {
13194     \exp_after:wN \__iow_wrap_fix_newline:w \l__iow_wrap_tl
13195     ^^J \q__iow_nil ^^J \s__iow_stop
13196   }
13197   \exp_after:wN \__iow_wrap_start:w \l__iow_wrap_tl
13198 }
13199 \cs_new:Npn \__iow_wrap_fix_newline:w #1 ^^J #2 ^^J
13200 {
13201   #1
13202   \if_meaning:w \q__iow_nil #2
13203   \__iow_use_i_delimit_by_s_stop:nw
13204   \fi:
13205   \c__iow_wrap_newline_marker_tl
13206   \__iow_wrap_fix_newline:w #2 ^^J
13207 }
13208 \cs_new_protected:Npn \__iow_wrap_start:w
13209 {
13210   \bool_set_false:N \l__iow_line_break_bool
13211   \tl_clear:N \l__iow_line_tl
13212   \tl_clear:N \l__iow_line_part_tl
13213   \tl_set:Nn \l__iow_wrap_tl { ~ \use_none:n }
13214   \int_zero:N \l__iow_indent_int
13215   \tl_clear:N \l__iow_indent_tl
13216   \__iow_wrap_chunk:nw { \l__iow_line_count_int }
13217 }

```

(End definition for `__iow_wrap_do:`, `__iow_wrap_fix_newline:w`, and `__iow_wrap_start:w`.)

`__iow_wrap_chunk:nw`
`__iow_wrap_next:nw`

The `chunk` and `next` auxiliaries are defined indirectly to obtain the expansions of `\c_catcode_other_space_tl` and `\c__iow_wrap_marker_tl` in their definition. The `next` auxiliary calls a function corresponding to the type of marker (its `##2`), which can be `newline` or `indent` or `unindent` or `end`. The first argument of the `chunk` auxiliary is a target number of characters and the second is some string to wrap. If the chunk is empty simply call `next`. Otherwise, set up a call to `__iow_wrap_line:nw`, including the indentation if the current line is empty, and including a trailing space (`#1`) before the `__iow_wrap_end_chunk:w` auxiliary.

```

13218 \cs_set_protected:Npn \__iow_tmp:w #1#2
13219 {
13220   \cs_new_protected:Npn \__iow_wrap_chunk:nw ##1##2 #2
13221   {
13222     \tl_if_empty:nTF {##2}
13223     {

```

```

13224         \tl_clear:N \l__iow_line_part_tl
13225         \__iow_wrap_next:nw {##1}
13226     }
13227     {
13228         \tl_if_empty:NTF \l__iow_line_tl
13229         {
13230             \__iow_wrap_line:nw
13231             { \l__iow_indent_tl }
13232             ##1 - \l__iow_indent_int ;
13233         }
13234         { \__iow_wrap_line:nw { } ##1 ; }
13235         ##2 #1
13236         \__iow_wrap_end_chunk:w 7 6 5 4 3 2 1 0 \s__iow_stop
13237     }
13238 }
13239 \cs_new_protected:Npn \__iow_wrap_next:nw ##1##2 #1
13240 { \use:c { __iow_wrap_##2:n } {##1} }
13241 }
13242 \exp_args:NVV \__iow_tmp:w \c_catcode_other_space_tl \c__iow_wrap_marker_tl

```

(End definition for `__iow_wrap_chunk:nw` and `__iow_wrap_next:nw`.)

```

\__iow_wrap_line:nw
\__iow_wrap_line_loop:w
\__iow_wrap_line_aux:Nw
\__iow_wrap_line_seven:nnnnnn
\__iow_wrap_line_end:NnnnnnnN
\__iow_wrap_line_end:nw
\__iow_wrap_end_chunk:w

```

This is followed by `{\langle string \rangle \langle intexpr \rangle ;}`. It stores the `\langle string \rangle` and up to `\langle intexpr \rangle` characters from the current chunk into `\l__iow_line_part_tl`. Characters are grabbed 8 at a time and left in `\l__iow_line_part_tl` by the `line_loop` auxiliary. When $k < 8$ remain to be found, the `line_aux` auxiliary calls the `line_end` auxiliary followed by (the single digit) k , then $7 - k$ empty brace groups, then the chunk's remaining characters. The `line_end` auxiliary leaves k characters from the chunk in the line part, then ends the assignment. Ignore the `\use_none:nnnnn` line for now. If the next character is a space the line can be broken there: store what we found into the result and get the next line. Otherwise some work is needed to find a break-point. So far we have ignored what happens if the chunk is shorter than the requested number of characters: this is dealt with by the `end_chunk` auxiliary, which gets treated like a character by the rest of the code. It ends up being called either as one of the arguments #2-#9 of the `line_loop` auxiliary or as one of the arguments #2-#8 of the `line_end` auxiliary. In both cases stop the assignment and work out how many characters are still needed. Notice that when we have exactly seven arguments to clean up, a `\exp_stop_f:` has to be inserted to stop the `\exp:w`. The weird `\use_none:nnnnn` ensures that the required data is in the right place.

```

13243 \cs_new_protected:Npn \__iow_wrap_line:nw #1
13244 {
13245     \tex_edef:D \l__iow_line_part_tl { \if_false: } \fi:
13246     #1
13247     \exp_after:wN \__iow_wrap_line_loop:w
13248     \int_value:w \int_eval:w
13249 }
13250 \cs_new:Npn \__iow_wrap_line_loop:w #1 ; #2#3#4#5#6#7#8#9
13251 {
13252     \if_int_compare:w #1 < 8 \exp_stop_f:
13253     \__iow_wrap_line_aux:Nw #1
13254     \fi:
13255     #2 #3 #4 #5 #6 #7 #8 #9
13256     \exp_after:wN \__iow_wrap_line_loop:w

```

```

13257     \int_value:w \int_eval:w #1 - 8 ;
13258   }
13259   \cs_new:Npn \__iow_wrap_line_aux:Nw #1#2#3 \exp_after:wN #4 ;
13260   {
13261     #2
13262     \exp_after:wN \__iow_wrap_line_end:NnnnnnnN
13263     \exp_after:wN #1
13264     \exp:w \exp_end_continue_f:w
13265     \exp_after:wN \exp_after:wN
13266     \if_case:w #1 \exp_stop_f:
13267       \prg_do_nothing:
13268     \or: \use_none:n
13269     \or: \use_none:nn
13270     \or: \use_none:nnn
13271     \or: \use_none:nnnn
13272     \or: \use_none:nnnnn
13273     \or: \use_none:nnnnnn
13274     \or: \__iow_wrap_line_seven:nnnnnnn
13275     \fi:
13276     { } { } { } { } { } { } { } { } { } #3
13277   }
13278   \cs_new:Npn \__iow_wrap_line_seven:nnnnnnn #1#2#3#4#5#6#7 { \exp_stop_f: }
13279   \cs_new:Npn \__iow_wrap_line_end:NnnnnnnN #1#2#3#4#5#6#7#8#9
13280   {
13281     #2 #3 #4 #5 #6 #7 #8
13282     \use_none:nnnnn \int_eval:w 8 - ; #9
13283     \token_if_eq_charcode:NNTF \c_space_token #9
13284     { \__iow_wrap_line_end:nw { } }
13285     { \if_false: { \fi: } \__iow_wrap_break:w #9 }
13286   }
13287   \cs_new:Npn \__iow_wrap_line_end:nw #1
13288   {
13289     \if_false: { \fi: }
13290     \__iow_wrap_store_do:n {#1}
13291     \__iow_wrap_next_line:w
13292   }
13293   \cs_new:Npn \__iow_wrap_end_chunk:w
13294     #1 \int_eval:w #2 - #3 ; #4#5 \s__iow_stop
13295   {
13296     \if_false: { \fi: }
13297     \exp_args:Nf \__iow_wrap_next:nw { \int_eval:n { #2 - #4 } }
13298   }

```

(End definition for __iow_wrap_line:nw and others.)

__iow_wrap_break:w Functions here are defined indirectly: __iow_tmp:w is eventually called with an “other”
 __iow_wrap_break_first:w space as its argument. The goal is to remove from \l__iow_line_part_tl the part
 __iow_wrap_break_none:w after the last space. In most cases this is done by repeatedly calling the **break_loop**
 __iow_wrap_break_loop:w auxiliary, which leaves “words” (delimited by spaces) until it hits the trailing space: then
 __iow_wrap_break_end:w its argument ##3 is ? __iow_wrap_break_end:w instead of a single token, and that
break_end auxiliary leaves in the assignment the line until the last space, then calls
 __iow_wrap_line_end:nw to finish up the line and move on to the next. If there is
 no space in \l__iow_line_part_tl then the **break_first** auxiliary calls the **break_**
none auxiliary. In that case, if the current line is empty, the complete word (including

##4, characters beyond what we had grabbed) is added to the line, making it over-long. Otherwise, the word is used for the following line (and the last space of the line so far is removed because it was inserted due to the presence of a marker).

```

13299 \cs_set_protected:Npn \__iow_tmp:w #1
13300 {
13301   \cs_new:Npn \__iow_wrap_break:w
13302   {
13303     \tex_edef:D \l__iow_line_part_tl
13304     { \if_false: } \fi:
13305     \exp_after:wN \__iow_wrap_break_first:w
13306     \l__iow_line_part_tl
13307     #1
13308     { ? \__iow_wrap_break_end:w }
13309     \s__iow_mark
13310   }
13311   \cs_new:Npn \__iow_wrap_break_first:w ##1 #1 ##2
13312   {
13313     \use_none:nn ##2 \__iow_wrap_break_none:w
13314     \__iow_wrap_break_loop:w ##1 #1 ##2
13315   }
13316   \cs_new:Npn \__iow_wrap_break_none:w ##1##2 #1 ##3 \s__iow_mark ##4 #1
13317   {
13318     \tl_if_empty:NTF \l__iow_line_tl
13319     { ##2 ##4 \__iow_wrap_line_end:nw { } }
13320     { \__iow_wrap_line_end:nw { \__iow_wrap_trim:N } ##2 ##4 #1 }
13321   }
13322   \cs_new:Npn \__iow_wrap_break_loop:w ##1 #1 ##2 #1 ##3
13323   {
13324     \use_none:n ##3
13325     ##1 #1
13326     \__iow_wrap_break_loop:w ##2 #1 ##3
13327   }
13328   \cs_new:Npn \__iow_wrap_break_end:w ##1 #1 ##2 ##3 #1 ##4 \s__iow_mark
13329   { ##1 \__iow_wrap_line_end:nw { } ##3 }
13330 }
13331 \exp_args:NV \__iow_tmp:w \c_catcode_other_space_tl

```

(End definition for __iow_wrap_break:w and others.)

__iow_wrap_next_line:w The special case where the end of a line coincides with the end of a chunk is detected here, to avoid a spurious empty line. Otherwise, call __iow_wrap_line:nw to find characters for the next line (remembering to account for the indentation).

```

13332 \cs_new_protected:Npn \__iow_wrap_next_line:w #1#2 \s__iow_stop
13333 {
13334   \tl_clear:N \l__iow_line_tl
13335   \token_if_eq_meaning:NNTF #1 \__iow_wrap_end_chunk:w
13336   {
13337     \tl_clear:N \l__iow_line_part_tl
13338     \bool_set_true:N \l__iow_line_break_bool
13339     \__iow_wrap_next:nw { \l__iow_line_target_int }
13340   }
13341   {
13342     \__iow_wrap_line:nw
13343     { \l__iow_indent_tl }

```

```

13344         \l__iow_line_target_int - \l__iow_indent_int ;
13345         #1 #2 \s__iow_stop
13346     }
13347 }

```

(End definition for __iow_wrap_next_line:w.)

__iow_wrap_allow_break:n This is called after a chunk has been wrapped. The \l__iow_line_part_tl typically ends with a space (except at the beginning of a line?), which we remove since the **allow-break** marker should not insert a space. Then move on with the next chunk, making sure to adjust the target number of characters for the line in case we did remove a space.

```

13348 \cs_new_protected:Npn \__iow_wrap_allow_break:n #1
13349 {
13350     \__kernel_tl_set:Nx \l__iow_line_tl
13351     { \l__iow_line_tl \__iow_wrap_trim:N \l__iow_line_part_tl }
13352     \bool_set_false:N \l__iow_line_break_bool
13353     \tl_if_empty:NTF \l__iow_line_part_tl
13354     { \__iow_wrap_chunk:nw {#1} }
13355     { \exp_args:Nf \__iow_wrap_chunk:nw { \int_eval:n { #1 + 1 } } }
13356 }

```

(End definition for __iow_wrap_allow_break:n.)

__iow_wrap_indent:n These functions are called after a chunk has been wrapped, when encountering **indent/unindent** markers. Add the line part (last line part of the previous chunk) to the line so far and reset a boolean denoting the presence of a line-break. Most importantly, add or remove one indent from the current indent (both the integer and the token list). Finally, continue wrapping.

__iow_wrap_unindent:n

```

13357 \cs_new_protected:Npn \__iow_wrap_indent:n #1
13358 {
13359     \tl_put_right:Nx \l__iow_line_tl { \l__iow_line_part_tl }
13360     \bool_set_false:N \l__iow_line_break_bool
13361     \int_add:Nn \l__iow_indent_int { \l__iow_one_indent_int }
13362     \tl_put_right:No \l__iow_indent_tl { \l__iow_one_indent_tl }
13363     \__iow_wrap_chunk:nw {#1}
13364 }
13365 \cs_new_protected:Npn \__iow_wrap_unindent:n #1
13366 {
13367     \tl_put_right:Nx \l__iow_line_tl { \l__iow_line_part_tl }
13368     \bool_set_false:N \l__iow_line_break_bool
13369     \int_sub:Nn \l__iow_indent_int { \l__iow_one_indent_int }
13370     \__kernel_tl_set:Nx \l__iow_indent_tl
13371     { \exp_after:wN \__iow_unindent:w \l__iow_indent_tl }
13372     \__iow_wrap_chunk:nw {#1}
13373 }

```

(End definition for __iow_wrap_indent:n and __iow_wrap_unindent:n.)

__iow_wrap_newline:n These functions are called after a chunk has been line-wrapped, when encountering a **newline/end** marker. Unless we just took a line-break, store the line part and the line so far into the whole \l__iow_wrap_tl, trimming a trailing space. In the **newline** case look for a new line (of length \l__iow_line_target_int) in a new chunk.

__iow_wrap_end:n

```

13374 \cs_new_protected:Npn \__iow_wrap_newline:n #1
13375 {

```

```

13376     \bool_if:NF \l__iow_line_break_bool
13377     { \__iow_wrap_store_do:n { \__iow_wrap_trim:N } }
13378     \bool_set_false:N \l__iow_line_break_bool
13379     \__iow_wrap_chunk:nw { \l__iow_line_target_int }
13380   }
13381   \cs_new_protected:Npn \__iow_wrap_end:n #1
13382   {
13383     \bool_if:NF \l__iow_line_break_bool
13384     { \__iow_wrap_store_do:n { \__iow_wrap_trim:N } }
13385     \bool_set_false:N \l__iow_line_break_bool
13386   }

```

(End definition for __iow_wrap_newline:n and __iow_wrap_end:n.)

__iow_wrap_store_do:n First add the last line part to the line, then append it to \l__iow_wrap_tl with the appropriate new line (with “run-on” text), possibly with its last space removed (#1 is empty or __iow_wrap_trim:N).

```

13387   \cs_new_protected:Npn \__iow_wrap_store_do:n #1
13388   {
13389     \__kernel_tl_set:Nx \l__iow_line_tl
13390     { \l__iow_line_tl \l__iow_line_part_tl }
13391     \__kernel_tl_set:Nx \l__iow_wrap_tl
13392     {
13393       \l__iow_wrap_tl
13394       \l__iow_newline_tl
13395       #1 \l__iow_line_tl
13396     }
13397     \tl_clear:N \l__iow_line_tl
13398   }

```

(End definition for __iow_wrap_store_do:n.)

__iow_wrap_trim:N Remove one trailing “other” space from the argument if present.

```

\__iow_wrap_trim:w
\__iow_wrap_trim_aux:w
13399   \cs_set_protected:Npn \__iow_tmp:w #1
13400   {
13401     \cs_new:Npn \__iow_wrap_trim:N ##1
13402     { \exp_after:wN \__iow_wrap_trim:w ##1 \s__iow_mark #1 \s__iow_mark \s__iow_stop }
13403     \cs_new:Npn \__iow_wrap_trim:w ##1 #1 \s__iow_mark
13404     { \__iow_wrap_trim_aux:w ##1 \s__iow_mark }
13405     \cs_new:Npn \__iow_wrap_trim_aux:w ##1 \s__iow_mark ##2 \s__iow_stop {##1}
13406   }
13407   \exp_args:NV \__iow_tmp:w \c_catcode_other_space_tl

```

(End definition for __iow_wrap_trim:N, __iow_wrap_trim:w, and __iow_wrap_trim_aux:w.)

13408 <@@=file>

20.4 File operations

\l__file_internal_tl Used as a short-term scratch variable.

```

13409   \tl_new:N \l__file_internal_tl

```

(End definition for \l__file_internal_tl.)

`\g_file_curr_dir_str` The name of the current file should be available at all times: the name itself is set dynamically.

`\g_file_curr_ext_str`

```

13410 \str_new:N \g_file_curr_dir_str
13411 \str_new:N \g_file_curr_ext_str
13412 \str_new:N \g_file_curr_name_str

```

(End definition for `\g_file_curr_dir_str`, `\g_file_curr_ext_str`, and `\g_file_curr_name_str`. These variables are documented on page 167.)

`\g__file_stack_seq` The input list of files is stored as a sequence stack. In package mode we can recover the information from the details held by L^AT_EX 2_ε (we must be in the preamble and loaded using `\usepackage` or `\RequirePackage`). As L^AT_EX 2_ε doesn't store directory and name separately, we stick to the same convention here. In pre-loading, `\@currnamestack` is empty so is skipped.

```

13413 \seq_new:N \g__file_stack_seq
13414 \group_begin:
13415   \cs_set_protected:Npn \__file_tmp:w #1#2#3
13416   {
13417     \tl_if_blank:nTF {#1}
13418     {
13419       \cs_set:Npn \__file_tmp:w ##1 " ##2 " ##3 \s__file_stop
13420       { { } {##2} { } }
13421       \seq_gput_right:Nx \g__file_stack_seq
13422       {
13423         \exp_after:wN \__file_tmp:w \tex_jobname:D
13424         " \tex_jobname:D " \s__file_stop
13425       }
13426     }
13427     {
13428       \seq_gput_right:Nn \g__file_stack_seq { { } {#1} {#2} }
13429       \__file_tmp:w
13430     }
13431   }
13432   \cs_if_exist:NT \@currnamestack
13433   {
13434     \tl_if_empty:NF \@currnamestack
13435     { \exp_after:wN \__file_tmp:w \@currnamestack }
13436   }
13437 \group_end:

```

(End definition for `\g__file_stack_seq`.)

`\g__file_record_seq` The total list of files used is recorded separately from the current file stack, as nothing is ever popped from this list. The current file name should be included in the file list! We will eventually copy the contents of `\@filelist`.

```

13438 \seq_new:N \g__file_record_seq

```

(End definition for `\g__file_record_seq`.)

`\l__file_base_name_tl` For storing the basename and full path whilst passing data internally.

`\l__file_full_name_tl`

```

13439 \tl_new:N \l__file_base_name_tl
13440 \tl_new:N \l__file_full_name_tl

```

(End definition for `\l__file_base_name_tl` and `\l__file_full_name_tl`.)

<code>\l__file_dir_str</code> <code>\l__file_ext_str</code> <code>\l__file_name_str</code>	<p>Used in parsing a path into parts: in contrast to the above, these are never used outside of the current module.</p> <pre> 13441 \str_new:N \l__file_dir_str 13442 \str_new:N \l__file_ext_str 13443 \str_new:N \l__file_name_str </pre> <p>(End definition for <code>\l__file_dir_str</code>, <code>\l__file_ext_str</code>, and <code>\l__file_name_str</code>.)</p>
<code>\l_file_search_path_seq</code>	<p>The current search path.</p> <pre> 13444 \seq_new:N \l_file_search_path_seq </pre> <p>(End definition for <code>\l_file_search_path_seq</code>. This variable is documented on page 168.)</p>
<code>\l__file_tmp_seq</code>	<p>Scratch space for comma list conversion.</p> <pre> 13445 \seq_new:N \l__file_tmp_seq </pre> <p>(End definition for <code>\l__file_tmp_seq</code>.)</p>
<h3>20.4.1 Internal auxiliaries</h3>	
<code>\s__file_stop</code>	<p>Internal scan marks.</p> <pre> 13446 \scan_new:N \s__file_stop </pre> <p>(End definition for <code>\s__file_stop</code>.)</p>
<code>\q__file_nil</code>	<p>Internal quarks.</p> <pre> 13447 \quark_new:N \q__file_nil </pre> <p>(End definition for <code>\q__file_nil</code>.)</p>
<code>__file_quark_if_nil_p:n</code> <code>__file_quark_if_nil:nTF</code>	<p>Branching quark conditional.</p> <pre> 13448 __kernel_quark_new_conditional:Nn __file_quark_if_nil:n { TF } </pre> <p>(End definition for <code>__file_quark_if_nil:nTF</code>.)</p>
<code>\q__file_recursion_tail</code> <code>\q__file_recursion_stop</code>	<p>Internal recursion quarks.</p> <pre> 13449 \quark_new:N \q__file_recursion_tail 13450 \quark_new:N \q__file_recursion_stop </pre> <p>(End definition for <code>\q__file_recursion_tail</code> and <code>\q__file_recursion_stop</code>.)</p>
<code>_file_if_recursion_tail_break:NN</code> <code>_file_if_recursion_tail_stop_do:Nn</code>	<p>Functions to query recursion quarks.</p> <pre> 13451 __kernel_quark_new_test:N _file_if_recursion_tail_stop:N 13452 __kernel_quark_new_test:N _file_if_recursion_tail_stop_do:nn </pre> <p>(End definition for <code>_file_if_recursion_tail_break:NN</code> and <code>_file_if_recursion_tail_stop_do:Nn</code>.)</p>

```

    \_kernel_file_name_sanitiz:n
    \_kernel_file_name_expand_loop:w
    \_kernel_file_name_expand_N_type:Nw
    \_kernel_file_name_expand_group:nw
    \_kernel_file_name_expand_space:w
    \_kernel_file_name_strip_quotes:n
    \_kernel_file_name_strip_quotes:nnw
    \_kernel_file_name_strip_quotes:nnn
    \_kernel_file_name_trim_spaces:n
    \_kernel_file_name_trim_spaces:nw
    \_kernel_file_name_trim_spaces_aux:n
    \_kernel_file_name_trim_spaces_aux:w

```

Expanding the file name without expanding active characters is done using the same token-by-token approach as for example case changing. The finale outcome only need be e-type expandable, so there is no need for the shuffling that is seen in other locations.

```

13453 \cs_new:Npn \_kernel_file_name_sanitiz:n #1
13454 {
13455   \exp_args:Ne \_kernel_file_name_trim_spaces:n
13456   {
13457     \exp_args:Ne \_kernel_file_name_strip_quotes:n
13458     {
13459       \_kernel_file_name_expand_loop:w #1
13460       \q__file_recursion_tail \q__file_recursion_stop
13461     }
13462   }
13463 }
13464 \cs_new:Npn \_kernel_file_name_expand_loop:w #1 \q__file_recursion_stop
13465 {
13466   \tl_if_head_is:N_type:nTF {#1}
13467   { \_kernel_file_name_expand_N_type:Nw }
13468   {
13469     \tl_if_head_is_group:nTF {#1}
13470     { \_kernel_file_name_expand_group:nw }
13471     { \_kernel_file_name_expand_space:w }
13472   }
13473   #1 \q__file_recursion_stop
13474 }
13475 \cs_new:Npn \_kernel_file_name_expand_N_type:Nw #1
13476 {
13477   \_file_if_recursion_tail_stop:N #1
13478   \bool_lazy_and:nnTF
13479   { \token_if_expandable_p:N #1 }
13480   {
13481     \bool_not_p:n
13482     {
13483       \bool_lazy_any_p:n
13484       {
13485         { \token_if_protected_macro_p:N #1 }
13486         { \token_if_protected_long_macro_p:N #1 }
13487         { \token_if_active_p:N #1 }
13488       }
13489     }
13490   }
13491   { \exp_after:wN \_kernel_file_name_expand_loop:w #1 }
13492   {
13493     \token_to_str:N #1
13494     \_kernel_file_name_expand_loop:w
13495   }
13496 }
13497 \cs_new:Npx \_kernel_file_name_expand_group:nw #1
13498 {
13499   \c_left_brace_str
13500   \exp_not:N \_kernel_file_name_expand_loop:w
13501   #1
13502   \c_right_brace_str
13503 }

```

```

13504 \exp_last_unbraced:NNo
13505 \cs_new:Npx \__kernel_file_name_expand_space:w \c_space_tl
13506 {
13507   \c_space_tl
13508   \exp_not:N \__kernel_file_name_expand_loop:w
13509 }

```

Quoting file name uses basically the same approach as for luaquotejobname: count the " tokens and remove them.

```

13510 \cs_new:Npn \__kernel_file_name_strip_quotes:n #1
13511 {
13512   \__kernel_file_name_strip_quotes:nnnw {#1} { 0 } { }
13513   #1 " \q_file_recursion_tail " \q_file_recursion_stop
13514 }
13515 \cs_new:Npn \__kernel_file_name_strip_quotes:nnnw #1#2#3#4 "
13516 {
13517   \__file_if_recursion_tail_stop_do:nn {#4}
13518   { \__kernel_file_name_strip_quotes:nnn {#1} {#2} {#3} }
13519   \__kernel_file_name_strip_quotes:nnnw {#1} { #2 + 1 } { #3#4 }
13520 }
13521 \cs_new:Npn \__kernel_file_name_strip_quotes:nnn #1#2#3
13522 {
13523   \int_if_even:nT {#2}
13524   {
13525     \__kernel_msg_expandable_error:nnn
13526     { kernel } { unbalanced-quote-in-filename } {#1}
13527   }
13528   #3
13529 }

```

Spaces need to be trimmed from the start of the name and from the end of any extension. However, the name we are passed might not have an extension: that means we have to look for one. If there is no extension, we still use the standard trimming function but deliberately prevent any spaces being removed at the end.

```

13530 \cs_new:Npn \__kernel_file_name_trim_spaces:n #1
13531 { \__kernel_file_name_trim_spaces:nw {#1} #1 . \q_file_nil . \s_file_stop }
13532 \cs_new:Npn \__kernel_file_name_trim_spaces:nw #1#2 . #3 . #4 \s_file_stop
13533 {
13534   \__file_quark_if_nil:nTF {#3}
13535   {
13536     \exp_args:Ne \__kernel_file_name_trim_spaces_aux:n
13537     { \tl_trim_spaces:n { #1 \s_file_stop } }
13538   }
13539   { \tl_trim_spaces:n {#1} }
13540 }
13541 \cs_new:Npn \__kernel_file_name_trim_spaces_aux:n #1
13542 { \__kernel_file_name_trim_spaces_aux:w #1 }
13543 \cs_new:Npn \__kernel_file_name_trim_spaces_aux:w #1 \s_file_stop {#1}

```

(End definition for __kernel_file_name_sanitize:n and others.)

```

\__kernel_file_name_quote:n
\__kernel_file_name_quote:nw
13544 \cs_new:Npn \__kernel_file_name_quote:n #1
13545 { \__kernel_file_name_quote:nw {#1} #1 ~ \q_file_nil \s_file_stop }
13546 \cs_new:Npn \__kernel_file_name_quote:nw #1 #2 ~ #3 \s_file_stop

```

```

13547 {
13548   \__file_quark_if_nil:nTF {#3}
13549   { #1 }
13550   { "#1" }
13551 }

```

(End definition for __kernel_file_name_quote:n and __kernel_file_name_quote:nw.)

\c__file_marker_tl The same idea as the marker for rescanning token lists: this pair of tokens cannot appear in a file that is being input.

```

13552 \tl_const:Nx \c__file_marker_tl { : \token_to_str:N : }

```

(End definition for \c__file_marker_tl.)

\file_get:nnTF The approach here is similar to that for \tl_set_rescan:Nnn. The file contents are grabbed as an argument delimited by \c__file_marker_tl. A few subtleties: braces in **\file_get:nnN** \if_false: ... \fi: to deal with possible alignment tabs, \tracingnesting to avoid **__file_get_aux:nnN** __file_get_do:Nw a warning about a group being closed inside the \scantokens, and \prg_return_true: is placed after the end-of-file marker.

```

13553 \cs_new_protected:Npn \file_get:nnN #1#2#3
13554 {
13555   \file_get:nnNF {#1} {#2} #3
13556   { \tl_set:Nn #3 { \q_no_value } }
13557 }
13558 \prg_new_protected_conditional:Npnn \file_get:nnN #1#2#3 { T , F , TF }
13559 {
13560   \file_get_full_name:nNTF {#1} \l__file_full_name_tl
13561   {
13562     \exp_args:NV \__file_get_aux:nnN
13563     \l__file_full_name_tl
13564     {#2} #3
13565     \prg_return_true:
13566   }
13567   { \prg_return_false: }
13568 }
13569 \cs_new_protected:Npx \__file_get_aux:nnN #1#2#3
13570 {
13571   \exp_not:N \if_false: { \exp_not:N \fi:
13572   \group_begin:
13573     \int_set_eq:NN \tex_tracingnesting:D \c_zero_int
13574     \exp_not:N \exp_args:No \tex_everyeof:D
13575     { \exp_not:N \c__file_marker_tl }
13576     #2 \scan_stop:
13577     \exp_not:N \exp_after:wN \exp_not:N \__file_get_do:Nw
13578     \exp_not:N \exp_after:wN #3
13579     \exp_not:N \exp_after:wN \exp_not:N \prg_do_nothing:
13580     \exp_not:N \tex_input:D
13581     \sys_if_engine_luatex:TF
13582     { {#1} }
13583     { \exp_not:N \__kernel_file_name_quote:n {#1} \scan_stop: }
13584   \exp_not:N \if_false: } \exp_not:N \fi:
13585 }
13586 \exp_args:Nno \use:nn
13587 { \cs_new_protected:Npn \__file_get_do:Nw #1#2 }

```

```

13588 { \c__file_marker_tl }
13589 {
13590   \group_end:
13591   \tl_set:No #1 {#2}
13592 }

```

(End definition for `\file_get:nnNTF` and others. These functions are documented on page 168.)

`__file_size:n` A copy of the primitive where it's available.

```

13593 \cs_new_eq:NN \__file_size:n \tex_filesize:D

```

(End definition for `__file_size:n`.)

`\file_full_name:n` File searching can be carried out if the `\pdffilesize` primitive or an equivalent is available. That of course means we need to arrange for everything else to here to be done by expansion too. We start off by sanitizing the name and quoting if required: we may need to remove those quotes, so the raw name is passed too.

```

\__file_full_name:n
\__file_full_name_aux:Nnn
\__file_full_name_aux:nN
\__file_name_cleanup:w
\__file_name_end:
\__file_name_ext_check:n
\__file_name_ext_check:nw
\__file_name_ext_check:nnw
\__file_name_ext_check:nn

```

```

13594 \cs_new:Npn \file_full_name:n #1
13595 {
13596   \exp_args:Ne \__file_full_name:n
13597   { \__kernel_file_name_sanitiz:n {#1} }
13598 }

```

First, we check if the file is just here: no mapping so we do not need the break part of the broader auxiliary. We are using the fact that the primitive here returns nothing if the file is entirely absent. For package mode, `\input@path` is a token list not a sequence.

```

13599 \cs_new:Npn \__file_full_name:n #1
13600 {
13601   \tl_if_blank:nF {#1}
13602   {
13603     \tl_if_blank:eTF { \__file_size:n {#1} }
13604     {
13605       \seq_map_tokens:Nn \l_file_search_path_seq
13606       { \__file_full_name_aux:Nnn \seq_map_break:n {#1} }
13607       \cs_if_exist:NT \input@path
13608       {
13609         \tl_map_tokens:Nn \input@path
13610         { \__file_full_name_aux:Nnn \tl_map_break:n {#1} }
13611       }
13612       \__file_name_end:
13613     }
13614     { \__file_ext_check:n {#1} }
13615   }
13616 }

```

Two pars to the auxiliary here so we can avoid doing quoting twice in the event we find the right file.

```

13617 \cs_new:Npn \__file_full_name_aux:Nnn #1#2#3
13618 { \exp_args:Ne \__file_full_name_aux:nN { \tl_to_str:n {#3} / #2 } #1 }
13619 \cs_new:Npn \__file_full_name_aux:nN #1 #2
13620 {
13621   \tl_if_blank:eF { \__file_size:n {#1} }
13622   {
13623     #2
13624     {

```

```

13625         \_file_ext_check:n {#1}
13626         \_file_name_cleanup:w
13627     }
13628 }
13629 }
13630 \cs_new:Npn \_file_name_cleanup:w #1 \_file_name_end: { }
13631 \cs_new:Npn \_file_name_end: { }

```

As T_EX automatically adds .tex if there is no extension, there is a little clean up to do here. First, make sure we are not in the directory part, saving that. Then check for an extension.

```

13632 \cs_new:Npn \_file_ext_check:n #1
13633 { \_file_ext_check:nw { / } #1 / \q_file_nil / \s_file_stop }
13634 \cs_new:Npn \_file_ext_check:nw #1 #2 / #3 / #4 \s_file_stop
13635 {
13636     \_file_quark_if_nil:nTF {#3}
13637     {
13638         \exp_args:No \_file_ext_check:nnw
13639         { \use_none:n #1 } {#2} #2 . \q_file_nil . \s_file_stop
13640     }
13641     { \_file_ext_check:nw { #1 #2 / } #3 / #4 \s_file_stop }
13642 }
13643 \cs_new:Npx \_file_ext_check:nnw #1#2#3 . #4 . #5 \s_file_stop
13644 {
13645     \exp_not:N \_file_quark_if_nil:nTF {#4}
13646     {
13647         \exp_not:N \_file_ext_check:nn
13648         { #1 #2 } { #1 #2 \tl_to_str:n { .tex } }
13649     }
13650     { #1 #2 }
13651 }
13652 \cs_new:Npn \_file_ext_check:nn #1#2
13653 {
13654     \tl_if_blank:eTF { \_file_size:n {#2} }
13655     {#1}
13656     {
13657         \int_compare:nNnTF
13658         { \_file_size:n {#1} } = { \_file_size:n {#2} }
13659         {#2}
13660         {#1}
13661     }
13662 }

```

Deal with the fact that the primitive might not be available.

```

13663 \cs_if_exist:NF \tex_filesize:D
13664 {
13665     \cs_gset:Npn \file_full_name:n #1
13666     {
13667         \_kernel_msg_expandable_error:nnn
13668         { kernel } { primitive-not-available }
13669         { \pdffilesize }
13670     }
13671 }
13672 \_kernel_msg_new:nnnn { kernel } { primitive-not-available }
13673 { Primitive~\token_to_str:N #1 not-available }

```

```

13674 {
13675     The-version-of-your-Tex-engine-does-not-provide-functionality-equivalent-to-
13676     the-#1-primitive.
13677 }

```

(End definition for `\file_full_name:n` and others. This function is documented on page 168.)

```

\file_get_full_name:nN
\file_get_full_name:VN
\file_get_full_name:nNTF
\file_get_full_name:VNTF
  \_file_get_full_name_search:nN

```

These functions pre-date using `\tex_filesize:D` for file searching, so are `get` functions with protection. To avoid having different search set ups, they are simply wrappers around the code above.

```

13678 \cs_new_protected:Npn \file_get_full_name:nN #1#2
13679 {
13680     \file_get_full_name:nNF {#1} #2
13681     { \tl_set:Nn #2 { \q_no_value } }
13682 }
13683 \cs_generate_variant:Nn \file_get_full_name:nN { V }
13684 \prg_new_protected_conditional:Npnn \file_get_full_name:nN #1#2 { T , F , TF }
13685 {
13686     \__kernel_tl_set:Nx #2
13687     { \file_full_name:n {#1} }
13688     \tl_if_empty:NTF #2
13689     { \prg_return_false: }
13690     { \prg_return_true: }
13691 }
13692 \cs_generate_variant:Nn \file_get_full_name:nNT { V }
13693 \cs_generate_variant:Nn \file_get_full_name:nNF { V }
13694 \cs_generate_variant:Nn \file_get_full_name:nNTF { V }

```

If `\tex_filesize:D` is not available, the way to test if a file exists is to try to open it: if it does not exist then `TeX` reports end-of-file. A search is made looking at each potential path in turn (starting from the current directory). The first location is of course treated as the correct one: this is done by jumping to `\prg_break_point:.` If nothing is found, `#2` is returned empty. A special case when there is no extension is that once the first location is found we test the existence of the file with `.tex` extension in that directory, and if it exists we include the `.tex` extension in the result.

```

13695 \cs_if_exist:NF \tex_filesize:D
13696 {
13697     \prg_set_protected_conditional:Npnn \file_get_full_name:nN #1#2 { T , F , TF }
13698     {
13699         \__kernel_tl_set:Nx \l__file_base_name_tl
13700         { \__kernel_file_name_sanitise:n {#1} }
13701         \__file_get_full_name_search:nN { } \use:n
13702         \seq_map_inline:Nn \l_file_search_path_seq
13703         { \__file_get_full_name_search:nN { ##1 / } \seq_map_break:n }
13704         \cs_if_exist:NT \input@path
13705         {
13706             \tl_map_inline:Nn \input@path
13707             { \__file_get_full_name_search:nN { ##1 } \tl_map_break:n }
13708         }
13709         \tl_set:Nn \l__file_full_name_tl { \q_no_value }
13710         \prg_break_point:
13711         \quark_if_no_value:NTF \l__file_full_name_tl
13712         {
13713             \ior_close:N \g__file_internal_ior

```



```

13714         \prg_return_false:
13715     }
13716     {
13717         \file_parse_full_name:VNNN \l__file_full_name_tl
13718         \l__file_dir_str \l__file_name_str \l__file_ext_str
13719         \str_if_empty:NT \l__file_ext_str
13720         {
13721             \__kernel_ior_open:No \g__file_internal_ior
13722             { \l__file_full_name_tl .tex }
13723             \ior_if_eof:NF \g__file_internal_ior
13724             { \tl_put_right:Nn \l__file_full_name_tl { .tex } }
13725         }
13726         \ior_close:N \g__file_internal_ior
13727         \tl_set_eq:NN #2 \l__file_full_name_tl
13728         \prg_return_true:
13729     }
13730 }
13731 }
13732 \cs_new_protected:Npn \__file_get_full_name_search:nN #1#2
13733 {
13734     \__kernel_tl_set:Nx \l__file_full_name_tl
13735     { \tl_to_str:n {#1} \l__file_base_name_tl }
13736     \__kernel_ior_open:No \g__file_internal_ior \l__file_full_name_tl
13737     \ior_if_eof:NF \g__file_internal_ior { #2 { \prg_break: } }
13738 }

```

(End definition for \file_get_full_name:nN, \file_get_full_name:nNTF, and __file_get_full_name_search:nN. These functions are documented on page 168.)

\g__file_internal_ior A reserved stream to test for file existence (if required), and for opening a shell.

```

13739 \ior_new:N \g__file_internal_ior

```

(End definition for \g__file_internal_ior.)

\file_md5five_hash:n Getting file details by expansion is relatively easy if a bit repetitive. As the MD5 function has a slightly different syntax from the other commands, there is a little cleaning up to do.

```

\file_size:n
\file_timestamp:n
\__file_details:nn
__file_details_aux:nn
__file_md5five_hash:n
13740 \cs_new:Npn \file_size:n #1
13741 { \__file_details:nn {#1} { size } }
13742 \cs_new:Npn \file_timestamp:n #1
13743 { \__file_details:nn {#1} { moddate } }
13744 \cs_new:Npn \__file_details:nn #1#2
13745 {
13746     \exp_args:Ne \__file_details_aux:nn
13747     { \file_full_name:n {#1} } {#2}
13748 }
13749 \cs_new:Npn \__file_details_aux:nn #1#2
13750 {
13751     \tl_if_blank:nF {#1}
13752     { \use:c { tex_file #2 :D } {#1} }
13753 }
13754 \cs_new:Npn \file_md5five_hash:n #1
13755 { \exp_args:Ne \__file_md5five_hash:n { \file_full_name:n {#1} } }
13756 \cs_new:Npn \__file_md5five_hash:n #1
13757 { \tex_md5fivesum:D file {#1} }

```

(End definition for `\file_md5ive_hash:n` and others. These functions are documented on page 170.)

`\file_hex_dump:nnn`

These are separate as they need multiple arguments *or* the file size. For LuaTeX, the emulation does not need the file size so we save a little on expansion.

`__file_hex_dump_auxi:nnn`

13758 `\cs_new:Npn \file_hex_dump:nnn #1#2#3`

13759 `{`

13760 `\exp_args:Neee __file_hex_dump_auxi:nnn`

13761 `{ \file_full_name:n {#1} }`

13762 `{ \int_eval:n {#2} }`

13763 `{ \int_eval:n {#3} }`

13764 `}`

13765 `\cs_new:Npn __file_hex_dump_auxi:nnn #1#2#3`

13766 `{`

13767 `\bool_lazy_any:nF`

13768 `{`

13769 `{ \tl_if_blank_p:n {#1} }`

13770 `{ \int_compare_p:nNn {#2} = 0 }`

13771 `{ \int_compare_p:nNn {#3} = 0 }`

13772 `}`

13773 `{`

13774 `\exp_args:Ne __file_hex_dump_auxii:nnnn`

13775 `{ __file_details_aux:nn {#1} { size } }`

13776 `{#1} {#2} {#3}`

13777 `}`

13778 `}`

13779 `\cs_new:Npn __file_hex_dump_auxii:nnnn #1#2#3#4`

13780 `{`

13781 `\int_compare:nNnTF {#3} > 0`

13782 `{ __file_hex_dump_auxiii:nnnn {#3} }`

13783 `{`

13784 `\exp_args:Ne __file_hex_dump_auxiii:nnnn`

13785 `{ \int_eval:n { #1 + #3 } }`

13786 `}`

13787 `{#1} {#2} {#4}`

13788 `}`

13789 `\cs_new:Npn __file_hex_dump_auxiii:nnnn #1#2#3#4`

13790 `{`

13791 `\int_compare:nNnTF {#4} > 0`

13792 `{ __file_hex_dump_auxiv:nnn {#4} }`

13793 `{`

13794 `\exp_args:Ne __file_hex_dump_auxiv:nnn`

13795 `{ \int_eval:n { #2 + #4 } }`

13796 `}`

13797 `{#1} {#3}`

13798 `}`

13799 `\cs_new:Npn __file_hex_dump_auxiv:nnn #1#2#3`

13800 `{`

13801 `\tex_filedump:D`

13802 `offset ~ \int_eval:n { #2 - 1 } ~`

13803 `length ~ \int_eval:n { #1 - #2 + 1 }`

13804 `{#3}`

13805 `}`

13806 `\cs_new:Npn \file_hex_dump:n #1`

13807 `{ \exp_args:Ne __file_hex_dump:n { \file_full_name:n {#1} } }`

```

13808 \sys_if_engine luatex:TF
13809 {
13810     \cs_new:Npn \__file_hex_dump:n #1
13811     {
13812         \tl_if_blank:nF {#1}
13813         { \tex_dump:D whole {#1} {#1} }
13814     }
13815 }
13816 {
13817     \cs_new:Npn \__file_hex_dump:n #1
13818     {
13819         \tl_if_blank:nF {#1}
13820         { \tex_dump:D length \tex_filesize:D {#1} {#1} }
13821     }
13822 }

```

(End definition for `\file_hex_dump:nnn` and others. These functions are documented on page 169.)

```

\file_get_hex_dump:nN
\file_get_hex_dump:nNTF
\file_get_md5fiver_hash:nN\file_get_size:nN
\file_get_md5fiver_hash:nN\file_get_size:nNTF
\file_get_timestamp:nN
\file_get_timestamp:nNTF
\__file_get_details:nnN
13823 \cs_new_protected:Npn \file_get_hex_dump:nN #1#2
13824 { \file_get_hex_dump:nNF {#1} #2 { \tl_set:Nn #2 { \q_no_value } } }
13825 \cs_new_protected:Npn \file_get_md5fiver_hash:nN #1#2
13826 { \file_get_md5fiver_hash:nNF {#1} #2 { \tl_set:Nn #2 { \q_no_value } } }
13827 \cs_new_protected:Npn \file_get_size:nN #1#2
13828 { \file_get_size:nNF {#1} #2 { \tl_set:Nn #2 { \q_no_value } } }
13829 \cs_new_protected:Npn \file_get_timestamp:nN #1#2
13830 { \file_get_timestamp:nNF {#1} #2 { \tl_set:Nn #2 { \q_no_value } } }
13831 \prg_new_protected_conditional:Npnn \file_get_hex_dump:nN #1#2 { T , F , TF }
13832 { \__file_get_details:nnN {#1} { hex_dump } #2 }
13833 \prg_new_protected_conditional:Npnn \file_get_md5fiver_hash:nN #1#2 { T , F , TF }
13834 { \__file_get_details:nnN {#1} { md5fiver_hash } #2 }
13835 \prg_new_protected_conditional:Npnn \file_get_size:nN #1#2 { T , F , TF }
13836 { \__file_get_details:nnN {#1} { size } #2 }
13837 \prg_new_protected_conditional:Npnn \file_get_timestamp:nN #1#2 { T , F , TF }
13838 { \__file_get_details:nnN {#1} { timestamp } #2 }
13839 \cs_new_protected:Npn \__file_get_details:nnN #1#2#3
13840 {
13841     \__kernel_tl_set:Nx #3
13842     { \use:c { file_ #2 :n } {#1} }
13843     \tl_if_empty:NTF #3
13844     { \prg_return_false: }
13845     { \prg_return_true: }
13846 }

```

Where the primitive is not available, issue an error: this is a little more conservative than absolutely needed, but does work.

```

13847 \cs_if_exist:NF \tex_filesize:D
13848 {
13849     \cs_set_protected:Npn \__file_get_details:nnN #1#2#3
13850     {
13851         \tl_clear:N #3
13852         \__kernel_msg_error:nnx
13853         { kernel } { primitive-not-available }
13854     }

```

```

13855         \token_to_str:N \(\pdf)file
13856         \str_case:nn {#2}
13857         {
13858             { hex_dump } { dump }
13859             { mdfive_hash } { mdfivesum }
13860             { timestamp } { moddate }
13861             { size } { size }
13862         }
13863     }
13864     \prg_return_false:
13865 }
13866 }

```

(End definition for \file_get_hex_dump:nNTF and others. These functions are documented on page 169.)

\file_get_hex_dump:nnnN
\file_get_hex_dump:nnnNTF

Custom code due to the additional arguments.

```

13867 \cs_new_protected:Npn \file_get_hex_dump:nnnN #1#2#3#4
13868 {
13869     \file_get_hex_dump:nnnNF {#1} {#2} {#3} #4
13870     { \tl_set:Nn #4 { \q_no_value } }
13871 }
13872 \prg_new_protected_conditional:Npnn \file_get_hex_dump:nnnN #1#2#3#4
13873 { T , F , TF }
13874 {
13875     \__kernel_tl_set:Nx #4
13876     { \file_hex_dump:nnn {#1} {#2} {#3} }
13877     \tl_if_empty:NTF #4
13878     { \prg_return_false: }
13879     { \prg_return_true: }
13880 }

```

(End definition for \file_get_hex_dump:nnnNTF. This function is documented on page 169.)

__file_str_cmp:nn

As we are doing a fixed-length “big” integer comparison, it is easiest to use the low-level behavior of string comparisons.

```

13881 \cs_new_eq:NN \__file_str_cmp:nn \tex_strcmp:D

```

(End definition for __file_str_cmp:nn.)

\file_compare_timestamp:p:nNn
\file_compare_timestamp:nNnTF
__file_compare_timestamp:nnN
__file_timestamp:n

Comparison of file date can be done by using the low-level nature of the string comparison functions.

```

13882 \prg_new_conditional:Npnn \file_compare_timestamp:nNn #1#2#3
13883 { p , T , F , TF }
13884 {
13885     \exp_args:Nee \__file_compare_timestamp:nnN
13886     { \file_full_name:n {#1} }
13887     { \file_full_name:n {#3} }
13888     #2
13889 }
13890 \cs_new:Npn \__file_compare_timestamp:nnN #1#2#3
13891 {
13892     \tl_if_blank:nTF {#1}
13893     {
13894         \if_charcode:w #3 <

```

```

13895         \prg_return_true:
13896     \else:
13897         \prg_return_false:
13898     \fi:
13899 }
13900 {
13901     \tl_if_blank:nTF {#2}
13902     {
13903         \if_charcode:w #3 >
13904         \prg_return_true:
13905         \else:
13906         \prg_return_false:
13907         \fi:
13908     }
13909     {
13910         \if_int_compare:w
13911         \__file_str_cmp:nn
13912         { \__file_timestamp:n {#1} }
13913         { \__file_timestamp:n {#2} }
13914         #3 0 \exp_stop_f:
13915         \prg_return_true:
13916         \else:
13917         \prg_return_false:
13918         \fi:
13919     }
13920 }
13921 }
13922 \cs_new_eq:NN \__file_timestamp:n \tex_filemoddate:D
13923 \cs_if_exist:NF \__file_timestamp:n
13924 {
13925     \prg_set_conditional:Npnn \file_compare_timestamp:nNn #1#2#3
13926     { p , T , F , TF }
13927     {
13928         \__kernel_msg_expandable_error:nnn
13929         { kernel } { primitive-not-available }
13930         { \ (pdf)filemoddate }
13931         \prg_return_false:
13932     }
13933 }

```

(End definition for `\file_compare_timestamp:nNnTF`, `__file_compare_timestamp:nnN`, and `__file_timestamp:n`. This function is documented on page 171.)

`\file_if_exist:nTF` The test for the existence of a file is a wrapper around the function to add a path to a file. If the file was found, the path contains something, whereas if the file was not located then the return value is empty.

```

13934 \prg_new_protected_conditional:Npnn \file_if_exist:n #1 { T , F , TF }
13935 {
13936     \file_get_full_name:nNTF {#1} \l__file_full_name_tl
13937     { \prg_return_true: }
13938     { \prg_return_false: }
13939 }

```

(End definition for `\file_if_exist:nTF`. This function is documented on page 168.)

`\file_if_exist_input:n` Input of a file with a test for existence. We do not define the T or TF variants because the most useful place to place the *<true code>* would be inconsistent with other conditionals.

```

13940 \cs_new_protected:Npn \file_if_exist_input:n #1
13941 {
13942   \file_get_full_name:nNT {#1} \l__file_full_name_tl
13943   { \__file_input:V \l__file_full_name_tl }
13944 }
13945 \cs_new_protected:Npn \file_if_exist_input:nF #1#2
13946 {
13947   \file_get_full_name:nNTF {#1} \l__file_full_name_tl
13948   { \__file_input:V \l__file_full_name_tl }
13949   {#2}
13950 }

```

(End definition for `\file_if_exist_input:n` and `\file_if_exist_input:nF`. These functions are documented on page 171.)

`\file_input_stop:` A simple rename.

```

13951 \cs_new_protected:Npn \file_input_stop: { \tex_endinput:D }

```

(End definition for `\file_input_stop:`. This function is documented on page 171.)

`__kernel_file_missing:n` An error message for a missing file, also used in `\ior_open:Nn`.

```

13952 \cs_new_protected:Npn \__kernel_file_missing:n #1
13953 {
13954   \__kernel_msg_error:nnx { kernel } { file-not-found }
13955   { \__kernel_file_name_sanitiz:n {#1} }
13956 }

```

(End definition for `__kernel_file_missing:n`.)

`\file_input:n` Loading a file is done in a safe way, checking first that the file exists and loading only if it does. Push the file name on the `\g__file_stack_seq`, and add it to the file list, either `\g__file_record_seq`, or `\@filelist` in package mode.

```

\__file_input:n
\__file_input:n
\__file_input:V
\__file_input_push:n
\__kernel_file_input_push:n
\__file_input_pop:
\__kernel_file_input_pop:
\__file_input_pop:nnn
13957 \cs_new_protected:Npn \file_input:n #1
13958 {
13959   \file_get_full_name:nNTF {#1} \l__file_full_name_tl
13960   { \__file_input:V \l__file_full_name_tl }
13961   { \__kernel_file_missing:n {#1} }
13962 }
13963 \cs_new_protected:Npx \__file_input:n #1
13964 {
13965   \exp_not:N \clist_if_exist:NTF \exp_not:N \@filelist
13966   { \exp_not:N \addtofilelist {#1} }
13967   { \seq_gput_right:Nn \exp_not:N \g__file_record_seq {#1} }
13968   \exp_not:N \__file_input_push:n {#1}
13969   \exp_not:N \tex_input:D
13970   \sys_if_engine luatex:TF
13971   { {#1} }
13972   { \exp_not:N \__kernel_file_name_quote:n {#1} \scan_stop: }
13973   \exp_not:N \__file_input_pop:
13974 }
13975 \cs_generate_variant:Nn \__file_input:n { V }

```

Keeping a track of the file data is easy enough: we store the separated parts so we do not need to parse them twice.

```

13976 \cs_new_protected:Npn \__file_input_push:n #1
13977 {
13978   \seq_gpush:Nx \g__file_stack_seq
13979   {
13980     { \g_file_curr_dir_str }
13981     { \g_file_curr_name_str }
13982     { \g_file_curr_ext_str }
13983   }
13984   \file_parse_full_name:nNNN {#1}
13985   \l__file_dir_str \l__file_name_str \l__file_ext_str
13986   \str_gset_eq:NN \g_file_curr_dir_str \l__file_dir_str
13987   \str_gset_eq:NN \g_file_curr_name_str \l__file_name_str
13988   \str_gset_eq:NN \g_file_curr_ext_str \l__file_ext_str
13989 }
13990 \cs_new_eq:NN \__kernel_file_input_push:n \__file_input_push:n
13991 \cs_new_protected:Npn \__file_input_pop:
13992 {
13993   \seq_gpop:NN \g__file_stack_seq \l__file_internal_tl
13994   \exp_after:wN \__file_input_pop:nnn \l__file_internal_tl
13995 }
13996 \cs_new_eq:NN \__kernel_file_input_pop: \__file_input_pop:
13997 \cs_new_protected:Npn \__file_input_pop:nnn #1#2#3
13998 {
13999   \str_gset:Nn \g_file_curr_dir_str {#1}
14000   \str_gset:Nn \g_file_curr_name_str {#2}
14001   \str_gset:Nn \g_file_curr_ext_str {#3}
14002 }

```

(End definition for `\file_input:n` and others. This function is documented on page 171.)

`\file_parse_full_name:n` The main parsing macro `\file_parse_full_name_apply:nN` passes the file name #1 through `__kernel_file_name_sanitiz:n` so that we have a single normalised way to treat files internally. `\file_parse_full_name:n` uses the former, with `\prg_do_nothing:` to leave each part of the name within a pair of braces.

```

14003 \cs_new:Npn \file_parse_full_name:n #1
14004 {
14005   \file_parse_full_name_apply:nN {#1}
14006   \prg_do_nothing:
14007 }
14008 \cs_new:Npn \file_parse_full_name_apply:nN #1
14009 {
14010   \exp_args:Ne \__file_parse_full_name_auxi:nN
14011   { \__kernel_file_name_sanitiz:n {#1} }
14012 }

```

`__file_parse_full_name_area:nw` splits the file name into chunks separated by /, until the last one is reached. The last chunk is the file name plus the extension, and everything before that is the path. When `__file_parse_full_name_area:nw` is done, it leaves the path within braces after the scan mark `\s__file_stop` and proceeds parsing the actual file name.

```

14013 \cs_new:Npn \__file_parse_full_name_auxi:nN #1

```

```

14014 {
14015     \__file_parse_full_name_area:nw { } #1
14016     / \s__file_stop
14017 }
14018 \cs_new:Npn \__file_parse_full_name_area:nw #1 #2 / #3 \s__file_stop
14019 {
14020     \tl_if_empty:nTF {#3}
14021     { \__file_parse_full_name_base:nw { } #2 . \s__file_stop {#1} }
14022     { \__file_parse_full_name_area:nw { #1 / #2 } #3 \s__file_stop }
14023 }

```

__file_parse_full_name_base:nw does roughly the same as above, but it separates the chunks at each period. However here there's some extra complications: In case #1 is empty, it is assumed that the extension is actually empty, and the file name is #2. Besides, an extra . has to be added to #2 because it is later removed in __file_parse_full_name_tidy:nnnN. In any case, if there's an extension, it is returned with a leading ..

__file_parse_full_name_base:nw

```

14024 \cs_new:Npn \__file_parse_full_name_base:nw #1 #2 . #3 \s__file_stop
14025 {
14026     \tl_if_empty:nTF {#3}
14027     {
14028         \tl_if_empty:nTF {#1}
14029         {
14030             \tl_if_empty:nTF {#2}
14031             { \__file_parse_full_name_tidy:nnnN { } { } }
14032             { \__file_parse_full_name_tidy:nnnN { .#2 } { } }
14033         }
14034         { \__file_parse_full_name_tidy:nnnN {#1} { .#2 } }
14035     }
14036     { \__file_parse_full_name_base:nw { #1 . #2 } #3 \s__file_stop }
14037 }

```

Now we just need to tidy some bits left loose before. The loop used in the two macros above start with a leading / and . in the file path an name, so here we need to remove them, except in the path, if it is a single /, in which case it's left as is. After all's done, pass to #4.

__file_parse_full_name_tidy:nnnN

```

14038 \cs_new:Npn \__file_parse_full_name_tidy:nnnN #1 #2 #3 #4
14039 {
14040     \exp_args:Nee #4
14041     {
14042         \str_if_eq:nnF {#3} { / } { \use_none:n }
14043         #3 \prg_do_nothing:
14044     }
14045     { \use_none:n #1 \prg_do_nothing: }
14046     {#2}
14047 }

```

(End definition for \file_parse_full_name:n and others. These functions are documented on page 169.)

\file_parse_full_name:nNNN

\file_parse_full_name:VNNN

```

14048 \cs_new_protected:Npn \file_parse_full_name:nNNN #1 #2 #3 #4
14049 {

```



```

14050 \file_parse_full_name_apply:nN {#1}
14051 \__file_full_name_assign:nnnNNN #2 #3 #4
14052 }
14053 \cs_new_protected:Npn \__file_full_name_assign:nnnNNN #1 #2 #3 #4 #5 #6
14054 {
14055 \str_set:Nn #4 {#1}
14056 \str_set:Nn #5 {#2}
14057 \str_set:Nn #6 {#3}
14058 }
14059 \cs_generate_variant:Nn \file_parse_full_name:nNNN { V }

```

(End definition for `\file_parse_full_name:nNNN`. This function is documented on page 169.)

`\file_show_list:` A function to list all files used to the log, without duplicates. In package mode, if
`\file_log_list:` `\@filelist` is still defined, we need to take this list of file names into account (we
`__file_list:N` capture it `\AtBeginDocument` into `\g__file_record_seq`), turning it to a string (this
`__file_list_aux:n` does not affect the commas of this comma list).

```

14060 \cs_new_protected:Npn \file_show_list: { \__file_list:N \msg_show:nnxxxx }
14061 \cs_new_protected:Npn \file_log_list: { \__file_list:N \msg_log:nnxxxx }
14062 \cs_new_protected:Npn \__file_list:N #1
14063 {
14064 \seq_clear:N \l__file_tmp_seq
14065 \clist_if_exist:NT \@filelist
14066 {
14067 \exp_args:NNx \seq_set_from_clist:Nn \l__file_tmp_seq
14068 { \tl_to_str:N \@filelist }
14069 }
14070 \seq_concat:NNN \l__file_tmp_seq \l__file_tmp_seq \g__file_record_seq
14071 \seq_remove_duplicates:N \l__file_tmp_seq
14072 #1 { LaTeX/kernel } { file-list }
14073 { \seq_map_function:NN \l__file_tmp_seq \__file_list_aux:n }
14074 { } { } { }
14075 }
14076 \cs_new:Npn \__file_list_aux:n #1 { \iow_newline: #1 }

```

(End definition for `\file_show_list:` and others. These functions are documented on page 171.)

When used as a package, there is a need to hold onto the standard file list as well as the new one here. File names recorded in `\@filelist` must be turned to strings before being added to `\g__file_record_seq`.

```

14077 \cs_if_exist:NT \@filelist
14078 {
14079 \AtBeginDocument
14080 {
14081 \exp_args:NNx \seq_set_from_clist:Nn \l__file_tmp_seq
14082 { \tl_to_str:N \@filelist }
14083 \seq_gconcat:NNN
14084 \g__file_record_seq
14085 \g__file_record_seq
14086 \l__file_tmp_seq
14087 }
14088 }

```

20.5 GetIdInfo

\GetIdInfo As documented in expl3.dtx this function extracts file name etc from an SVN Id line. This used to be how we got version number and so on in all modules, so it had to be defined in l3bootstrap. Now it's more convenient to define it after we have set up quite a lot of tools, and l3file seems the least unreasonable place for it.

The idea here is to extract out the information needed from a standard SVN Id line, but to avoid a line that would get changed when the file is checked in. Hence the fact that none of the lines here include both a dollar sign and the Id keyword!

```

14089 \cs_new_protected:Npn \GetIdInfo
14090 {
14091   \tl_clear_new:N \ExplFileDescription
14092   \tl_clear_new:N \ExplFileDate
14093   \tl_clear_new:N \ExplFileName
14094   \tl_clear_new:N \ExplFileExtension
14095   \tl_clear_new:N \ExplFileVersion
14096   \group_begin:
14097   \char_set_catcode_space:n { 32 }
14098   \exp_after:wN
14099   \group_end:
14100   \__file_id_info_auxi:w
14101 }

```

A first check for a completely empty SVN field. If that is not the case, there is a second case when a file created using svn cp but has not been checked in. That leaves a special marker -1 version, which has no further data. Dealing correctly with that is the reason for the space in the line to use __file_id_info_auxii:w.

```

14102 \cs_new_protected:Npn \__file_id_info_auxi:w $ #1 $ #2
14103 {
14104   \tl_set:Nn \ExplFileDescription {#2}
14105   \str_if_eq:nnTF {#1} { Id }
14106   {
14107     \tl_set:Nn \ExplFileDate { 0000/00/00 }
14108     \tl_set:Nn \ExplFileName { [unknown] }
14109     \tl_set:Nn \ExplFileExtension { [unknown~extension] }
14110     \tl_set:Nn \ExplFileVersion {-1}
14111   }
14112   { \__file_id_info_auxii:w #1 ~ \s__file_stop }
14113 }

```

Here, #1 is Id, #2 is the file name, #3 is the extension, #4 is the version, #5 is the check in date and #6 is the check in time and user, plus some trailing spaces. If #4 is the marker -1 value then #5 and #6 are empty.

```

14114 \cs_new_protected:Npn \__file_id_info_auxii:w
14115   #1 ~ #2.#3 ~ #4 ~ #5 ~ #6 \s__file_stop
14116 {
14117   \tl_set:Nn \ExplFileName {#2}
14118   \tl_set:Nn \ExplFileExtension {#3}
14119   \tl_set:Nn \ExplFileVersion {#4}
14120   \str_if_eq:nnTF {#4} {-1}
14121   { \tl_set:Nn \ExplFileDate { 0000/00/00 } }
14122   { \__file_id_info_auxiii:w #5 - 0 - 0 - \s__file_stop }
14123 }

```

Convert an SVN-style date into a L^AT_EX-style one.

```

14124 \cs_new_protected:Npn \__file_id_info_auxiii:w #1 - #2 - #3 - #4 \s__file_stop
14125 { \tl_set:Nn \ExplFileDate { #1/#2/#3 } }

```

(End definition for \GetIdInfo and others. This function is documented on page 7.)

20.6 Checking the version of kernel dependencies

```

\__kernel_dependency_version_check:Nn
\__kernel_dependency_version_check:nn
\__file_kernel_dependency_compare:nnn
\__file_parse_version:w

```

This function is responsible for checking if dependencies of the L^AT_EX3 kernel match the version preloaded in the L^AT_EX 2_ε kernel. If versions don't match, the function attempts to tell why by searching for a possible stray format file.

The function starts by checking that the kernel date is defined, and if not zero is used to force the error route. The kernel date is then compared with the argument requested date (usually the packaging date of the dependency). If the kernel date is less than the required date, it's an error and the loading should abort.

```

14126 \cs_new_protected:Npn \__kernel_dependency_version_check:Nn #1
14127 { \exp_args:NV \__kernel_dependency_version_check:nn #1 }
14128 \cs_new_protected:Npn \__kernel_dependency_version_check:nn #1
14129 {
14130   \cs_if_exist:NTF \c__kernel_expl_date_tl
14131   {
14132     \exp_args:NV \__file_kernel_dependency_compare:nnn
14133     \c__kernel_expl_date_tl {#1}
14134   }
14135   { \__file_kernel_dependency_compare:nnn { 0000-00-00 } {#1} }
14136 }
14137 \cs_new_protected:Npn \__file_kernel_dependency_compare:nnn #1 #2 #3
14138 {
14139   \int_compare:nNnT
14140   { \__file_parse_version:w #1 \s__file_stop } <
14141   { \__file_parse_version:w #2 \s__file_stop }
14142   { \__file_mismatched_dependency_error:nn {#2} {#3} }
14143 }
14144 \cs_new:Npn \__file_parse_version:w #1 - #2 - #3 \s__file_stop {#1#2#3}

```

If the versions differ, then we try to give the user some guidance. This function starts by taking the engine name \c_sys_engine_str and replacing tex by latex, then building a command of the form: kpsewhich -all -engine=<engine> <format>[-dev].fmt to query the format files available. A shell is opened and each line is read into a sequence.

```

\__file_mismatched_dependency_error:nn

```

```

14145 \cs_new_protected:Npn \__file_mismatched_dependency_error:nn #1 #2
14146 {
14147   \exp_args:NNx \ior_shell_open:Nn \g__file_internal_ior
14148   {
14149     kpsewhich ~ --all ~
14150     --engine = \c_sys_engine_exec_str
14151     \c_space_tl \c_sys_engine_format_str
14152     \bool_lazy_and:nnF
14153     { \tl_if_exist_p:N \development@branch@name }
14154     { ! \tl_if_empty_p:N \development@branch@name }
14155     { -dev } .fmt
14156   }
14157   \seq_clear:N \l__file_tmp_seq
14158   \ior_map_inline:Nn \g__file_internal_ior

```

```

14159     { \seq_put_right:Nn \l__file_tmp_seq {##1} }
14160 \ior_close:N \g__file_internal_ior
14161 \__kernel_msg_error:nnnn { kernel } { mismatched-support-file }
14162     {#1} {#2}

```

And finish by ending the current file.

```

14163 \tex_endinput:D
14164 }
14165 % \begin{macrocode}
14166 %
14167 % Now define the actual error message:
14168 % \begin{macrocode}
14169 \__kernel_msg_new:nnnn { kernel } { mismatched-support-file }
14170 {
14171     Mismatched~LaTeX~support~files~detected. \\
14172     Loading~'~#2'~aborted!

```

`\c__kernel_expl_date_tl` may not exist, due to an older format, so only print the dates when the sentinel token list exists:

```

14173 \tl_if_exist:NT \c__kernel_expl_date_tl
14174 {
14175     \\ \\
14176     The~L3~programming~layer~in~the~LaTeX~format \\
14177     is~dated~\c__kernel_expl_date_tl,~but~in~your~TeX~
14178     tree~the~files~require \\ at~least~#1.
14179 }
14180 }
14181 {

```

The sequence containing the format files should have exactly one item: the format file currently being run. If that's the case, the cause of the error is not that, so print a generic help with some possible causes. If more than one format file was found, then print the list to the user, with appropriate indications of what's in the system and what's in the user tree.

```

14182 \int_compare:nNnTF { \seq_count:N \l__file_tmp_seq } > 1
14183 {
14184     The~cause~seems~to~be~an~old~format~file~in~the~user~tree. \\
14185     LaTeX~found~these~files:
14186     \seq_map_tokens:Nn \l__file_tmp_seq { \\~~~\use:n } \\
14187     Try~deleting~the~file~in~the~user~tree~then~run~LaTeX~again.
14188 }
14189 {
14190     The~most~likely~causes~are:
14191     \\~~~A~recent~format~generation~failed;
14192     \\~~~A~stray~format~file~in~the~user~tree~which~needs~
14193     to~be~removed~or~rebuilt;
14194     \\~~~You~are~running~a~manually~installed~version~of~#2 \\
14195     \ \ \ which~is~incompatible~with~the~version~in~LaTeX. \\
14196 }
14197 \\
14198 LaTeX~will~abort~loading~the~incompatible~support~files~
14199 but~this~may~lead~to \\ later~errors.~Please~ensure~that~
14200 your~LaTeX~format~is~correctly~regenerated.
14201 }

```

(End definition for `__kernel_dependency_version_check:Nn` and others.)

20.7 Messages

```

14202 \__kernel_msg_new:nnnn { kernel } { file-not-found }
14203 { File~'#1'~not-found. }
14204 {
14205     The~requested~file~could~not~be~found~in~the~current~directory,~
14206     in~the~TeX~search~path~or~in~the~LaTeX~search~path.
14207 }
14208 \__kernel_msg_new:nnn { kernel } { file-list }
14209 {
14210     >~File~List~<
14211     #1 \\
14212     .....
14213 }
14214 \__kernel_msg_new:nnnn { kernel } { unbalanced-quote-in-filename }
14215 { Unbalanced~quotes~in~file~name~'#1'. }
14216 {
14217     File~names~must~contain~balanced~numbers~of~quotes~(").
14218 }
14219 \__kernel_msg_new:nnnn { kernel } { iow-indent }
14220 { Only~#1 (arg~1)~allows~#2 }
14221 {
14222     The~command~#2 can~only~be~used~in~messages~
14223     which~will~be~wrapped~using~#1.
14224     \tl_if_empty:nF {#3} { ~ It~was~called~with~argument~'#3'. }
14225 }

```

20.8 Functions delayed from earlier modules

<@@=sys>

\c_sys_platform_str Detecting the platform on LuaTeX is easy: for other engines, we use the fact that the two common cases have special null files. It is possible to probe further (see package `platform`), but that requires shell escape and seems unlikely to be useful. This is set up here as it requires file searching.

```

14226 \sys_if_engine luatex:TF
14227 {
14228     \str_const:Nx \c_sys_platform_str
14229     { \tex_directlua:D { tex.print(os.type) } }
14230 }
14231 {
14232     \file_if_exist:nTF { nul: }
14233     {
14234         \file_if_exist:nF { /dev/null }
14235         { \str_const:Nn \c_sys_platform_str { windows } }
14236     }
14237     {
14238         \file_if_exist:nT { /dev/null }
14239         { \str_const:Nn \c_sys_platform_str { unix } }
14240     }
14241 }
14242 \cs_if_exist:NF \c_sys_platform_str
14243 { \str_const:Nn \c_sys_platform_str { unknown } }

```

(End definition for `\c_sys_platform_str`. This variable is documented on page [116](#).)

```

\sys_if_platform_unix_p: We can now set up the tests.
\sys_if_platform_unix:TF 14244 \clist_map_inline:nn { unix , windows }
\sys_if_platform_windows_p: 14245 {
\sys_if_platform_windows:TF 14246   \__file_const:nn { sys_if_platform_ #1 }
                             14247   { \str_if_eq_p:Vn \c_sys_platform_str { #1 } }
                             14248 }

(End definition for \sys_if_platform_unix:TF and \sys_if_platform_windows:TF. These functions are
documented on page 116.)

14249 \</package>

```

21 l3skip implementation

```

14250 \*package>
14251 \@@=dim>

```

21.1 Length primitives renamed

```

\if_dim:w Primitives renamed.
\__dim_eval:w 14252 \cs_new_eq:NN \if_dim:w \tex_ifdim:D
\__dim_eval_end: 14253 \cs_new_eq:NN \__dim_eval:w \tex_dimexpr:D
                  14254 \cs_new_eq:NN \__dim_eval_end: \tex_relax:D

(End definition for \if_dim:w, \__dim_eval:w, and \__dim_eval_end:. This function is documented on
page 186.)

```

21.2 Internal auxiliaries

```

\s__dim_mark Internal scan marks.
\s__dim_stop 14255 \scan_new:N \s__dim_mark
              14256 \scan_new:N \s__dim_stop

(End definition for \s__dim_mark and \s__dim_stop.)

\__dim_use_none_delimit_by_s_stop:w Functions to gobble up to a scan mark.
14257 \cs_new:Npn \__dim_use_none_delimit_by_s_stop:w #1 \s__dim_stop { }

(End definition for \__dim_use_none_delimit_by_s_stop:w.)

```

21.3 Creating and initialising dim variables

```

\dim_new:N Allocating <dim> registers ...
\dim_new:c 14258 \cs_new_protected:Npn \dim_new:N #1
           14259 {
           14260   \__kernel_chk_if_free_cs:N #1
           14261   \cs:w newdimen \cs_end: #1
           14262 }
           14263 \cs_generate_variant:Nn \dim_new:N { c }

(End definition for \dim_new:N. This function is documented on page 172.)

```

\dim_const:Nn Contrarily to integer constants, we cannot avoid using a register, even for constants. We cannot use **\dim_gset:Nn** because debugging code would complain that the constant is not a global variable. Since **\dim_const:Nn** does not need to be fast, use **\dim_eval:n** to avoid needing a debugging patch that wraps the expression in checking code.

```

14264 \cs_new_protected:Npn \dim_const:Nn #1#2
14265 {
14266   \dim_new:N #1
14267   \tex_global:D #1 ~ \dim_eval:n {#2} \scan_stop:
14268 }
14269 \cs_generate_variant:Nn \dim_const:Nn { c }

```

(End definition for **\dim_const:Nn**. This function is documented on page 172.)

\dim_zero:N Reset the register to zero. Using **\c_zero_skip** deals with the case where the variable passed is incorrectly a skip (for example a $\text{L}^{\text{A}}\text{T}_{\text{E}}\text{X} 2_{\epsilon}$ length).

```

\dim_zero:c
\dim_gzero:N
\dim_gzero:c
14270 \cs_new_protected:Npn \dim_zero:N #1 { #1 \c_zero_skip }
14271 \cs_new_protected:Npn \dim_gzero:N #1
14272 { \tex_global:D #1 \c_zero_skip }
14273 \cs_generate_variant:Nn \dim_zero:N { c }
14274 \cs_generate_variant:Nn \dim_gzero:N { c }

```

(End definition for **\dim_zero:N** and **\dim_gzero:N**. These functions are documented on page 172.)

\dim_zero_new:N Create a register if needed, otherwise clear it.

```

\dim_zero_new:c
\dim_gzero_new:N
\dim_gzero_new:c
14275 \cs_new_protected:Npn \dim_zero_new:N #1
14276 { \dim_if_exist:NTF #1 { \dim_zero:N #1 } { \dim_new:N #1 } }
14277 \cs_new_protected:Npn \dim_gzero_new:N #1
14278 { \dim_if_exist:NTF #1 { \dim_gzero:N #1 } { \dim_new:N #1 } }
14279 \cs_generate_variant:Nn \dim_zero_new:N { c }
14280 \cs_generate_variant:Nn \dim_gzero_new:N { c }

```

(End definition for **\dim_zero_new:N** and **\dim_gzero_new:N**. These functions are documented on page 172.)

\dim_if_exist_p:N Copies of the **cs** functions defined in **l3basics**.

```

\dim_if_exist_p:c \prg_new_eq_conditional:NNn \dim_if_exist:N \cs_if_exist:N
\dim_if_exist:NTF { TF , T , F , p }
\dim_if_exist:cTF \prg_new_eq_conditional:NNn \dim_if_exist:c \cs_if_exist:c
14283
14284 { TF , T , F , p }

```

(End definition for **\dim_if_exist:NTF**. This function is documented on page 172.)

21.4 Setting dim variables

\dim_set:Nn Setting dimensions is easy enough but when debugging we want both to check that the variable is correctly local/global and to wrap the expression in some code. The **\scan_stop:** deals with the case where the variable passed is a skip (for example a $\text{L}^{\text{A}}\text{T}_{\text{E}}\text{X} 2_{\epsilon}$ length).

```

\dim_set:c
\dim_gset:Nn
\dim_gset:c
14285 \cs_new_protected:Npn \dim_set:Nn #1#2
14286 { #1 ~ \__dim_eval:w #2 \__dim_eval_end: \scan_stop: }
14287 \cs_new_protected:Npn \dim_gset:Nn #1#2
14288 { \tex_global:D #1 ~ \__dim_eval:w #2 \__dim_eval_end: \scan_stop: }
14289 \cs_generate_variant:Nn \dim_set:Nn { c }
14290 \cs_generate_variant:Nn \dim_gset:Nn { c }

```

(End definition for `\dim_set:Nn` and `\dim_gset:Nn`. These functions are documented on page 173.)

```

\dim_set_eq:NN All straightforward, with a \scan_stop: to deal with the case where #1 is (incorrectly)
\dim_set_eq:cN a skip.
\dim_set_eq:Nc
\dim_set_eq:cc 14291 \cs_new_protected:Npn \dim_set_eq:NN #1#2
14292 { #1 = #2 \scan_stop: }
\dim_gset_eq:NN 14293 \cs_generate_variant:Nn \dim_set_eq:NN { c , Nc , cc }
\dim_gset_eq:cN 14294 \cs_new_protected:Npn \dim_gset_eq:NN #1#2
\dim_gset_eq:Nc 14295 { \tex_global:D #1 = #2 \scan_stop: }
\dim_gset_eq:cc 14296 \cs_generate_variant:Nn \dim_gset_eq:NN { c , Nc , cc }

```

(End definition for `\dim_set_eq:NN` and `\dim_gset_eq:NN`. These functions are documented on page 173.)

```

\dim_add:Nn Using by here deals with the (incorrect) case \dimen123. Using \scan_stop: deals with
\dim_add:cN skip variables. Since debugging checks that the variable is correctly local/global, the
\dim_gadd:Nn global versions cannot be defined as \tex_global:D followed by the local versions. The
\dim_gadd:cN debugging code is inserted by \__dim_tmp:w.
\dim_sub:Nn 14297 \cs_new_protected:Npn \dim_add:Nn #1#2
\dim_sub:cN 14298 { \tex_advance:D #1 by \__dim_eval:w #2 \__dim_eval_end: \scan_stop: }
\dim_gsub:Nn 14299 \cs_new_protected:Npn \dim_gadd:Nn #1#2
\dim_gsub:cN 14300 {
14301 \tex_global:D \tex_advance:D #1 by
14302 \__dim_eval:w #2 \__dim_eval_end: \scan_stop:
14303 }
14304 \cs_generate_variant:Nn \dim_add:Nn { c }
14305 \cs_generate_variant:Nn \dim_gadd:Nn { c }
14306 \cs_new_protected:Npn \dim_sub:Nn #1#2
14307 { \tex_advance:D #1 by - \__dim_eval:w #2 \__dim_eval_end: \scan_stop: }
14308 \cs_new_protected:Npn \dim_gsub:Nn #1#2
14309 {
14310 \tex_global:D \tex_advance:D #1 by
14311 -\__dim_eval:w #2 \__dim_eval_end: \scan_stop:
14312 }
14313 \cs_generate_variant:Nn \dim_sub:Nn { c }
14314 \cs_generate_variant:Nn \dim_gsub:Nn { c }

```

(End definition for `\dim_add:Nn` and others. These functions are documented on page 173.)

21.5 Utilities for dimension calculations

```

\dim_abs:n Functions for min, max, and absolute value with only one evaluation. The absolute value
\__dim_abs:N is evaluated by removing a leading - if present.
\dim_max:nn 14315 \cs_new:Npn \dim_abs:n #1
\dim_min:nn 14316 {
\__dim_maxmin:wwN 14317 \exp_after:wN \__dim_abs:N
14318 \dim_use:N \__dim_eval:w #1 \__dim_eval_end:
14319 }
14320 \cs_new:Npn \__dim_abs:N #1
14321 { \if_meaning:w - #1 \else: \exp_after:wN #1 \fi: }
14322 \cs_new:Npn \dim_max:nn #1#2
14323 {
14324 \dim_use:N \__dim_eval:w \exp_after:wN \__dim_maxmin:wwN
14325 \dim_use:N \__dim_eval:w #1 \exp_after:wN ;

```



```

14326     \dim_use:N \__dim_eval:w #2 ;
14327     >
14328     \__dim_eval_end:
14329   }
14330 \cs_new:Npn \dim_min:nn #1#2
14331 {
14332   \dim_use:N \__dim_eval:w \exp_after:wN \__dim_maxmin:wwN
14333   \dim_use:N \__dim_eval:w #1 \exp_after:wN ;
14334   \dim_use:N \__dim_eval:w #2 ;
14335   <
14336   \__dim_eval_end:
14337 }
14338 \cs_new:Npn \__dim_maxmin:wwN #1 ; #2 ; #3
14339 {
14340   \if_dim:w #1 #3 #2 ~
14341     #1
14342   \else:
14343     #2
14344   \fi:
14345 }

```

(End definition for `\dim_abs:n` and others. These functions are documented on page 173.)

`\dim_ratio:nn` With dimension expressions, something like `10 pt * (5 pt / 10 pt)` does not work. Instead, the ratio part needs to be converted to an integer expression. Using `\int_value:w` forces everything into `sp`, avoiding any decimal parts.

```

14346 \cs_new:Npn \dim_ratio:nn #1#2
14347 { \__dim_ratio:n {#1} / \__dim_ratio:n {#2} }
14348 \cs_new:Npn \__dim_ratio:n #1
14349 { \int_value:w \__dim_eval:w (#1) \__dim_eval_end: }

```

(End definition for `\dim_ratio:nn` and `__dim_ratio:n`. This function is documented on page 174.)

21.6 Dimension expression conditionals

`\dim_compare_p:nNn` Simple comparison.

`\dim_compare:nNnTF`

```

14350 \prg_new_conditional:Npnn \dim_compare:nNn #1#2#3 { p , T , F , TF }
14351 {
14352   \if_dim:w \__dim_eval:w #1 #2 \__dim_eval:w #3 \__dim_eval_end:
14353   \prg_return_true: \else: \prg_return_false: \fi:
14354 }

```

(End definition for `\dim_compare:nNnTF`. This function is documented on page 174.)

`\dim_compare_p:n` This code is adapted from the `\int_compare:nTF` function. First make sure that there is at least one relation operator, by evaluating a dimension expression with a trailing `__dim_compare_error:`. Just like for integers, the looping auxiliary `__dim_compare:wNN` closes a primitive conditional and opens a new one. It is actually easier to grab a dimension operand than an integer one, because once evaluated, dimensions all end with `pt` (with category other). Thus we do not need specific auxiliaries for the three “simple” relations `<`, `=`, and `>`.

`\dim_compare:nTF`

```

\__dim_compare:w
\__dim_compare:wNN
\__dim_compare:=:w
\__dim_compare!:w
\__dim_compare<:w
\__dim_compare>:w
\__dim_compare_error:
14355 \prg_new_conditional:Npnn \dim_compare:n #1 { p , T , F , TF }
14356 {
14357   \exp_after:wN \__dim_compare:w

```

```

14358     \dim_use:N \__dim_eval:w #1 \__dim_compare_error:
14359   }
14360 \cs_new:Npn \__dim_compare:w #1 \__dim_compare_error:
14361 {
14362   \exp_after:wN \if_false: \exp:w \exp_end_continue_f:w
14363   \__dim_compare:wNN #1 ? { = \__dim_compare_end:w \else: } \s__dim_stop
14364 }
14365 \exp_args:Nno \use:nn
14366 { \cs_new:Npn \__dim_compare:wNN #1 } { \tl_to_str:n {pt} #2#3 }
14367 {
14368   \if_meaning:w = #3
14369   \use:c { __dim_compare_#2:w }
14370   \fi:
14371   #1 pt \exp_stop_f:
14372   \prg_return_false:
14373   \exp_after:wN \__dim_use_none_delimit_by_s_stop:w
14374   \fi:
14375   \reverse_if:N \if_dim:w #1 pt #2
14376   \exp_after:wN \__dim_compare:wNN
14377   \dim_use:N \__dim_eval:w #3
14378 }
14379 \cs_new:cpn { __dim_compare_! :w }
14380 #1 \reverse_if:N #2 ! #3 = { #1 #2 = #3 }
14381 \cs_new:cpn { __dim_compare_ = :w }
14382 #1 \__dim_eval:w = { #1 \__dim_eval:w }
14383 \cs_new:cpn { __dim_compare_ < :w }
14384 #1 \reverse_if:N #2 < #3 = { #1 #2 > #3 }
14385 \cs_new:cpn { __dim_compare_ > :w }
14386 #1 \reverse_if:N #2 > #3 = { #1 #2 < #3 }
14387 \cs_new:Npn \__dim_compare_end:w #1 \prg_return_false: #2 \s__dim_stop
14388 { #1 \prg_return_false: \else: \prg_return_true: \fi: }
14389 \cs_new_protected:Npn \__dim_compare_error:
14390 {
14391   \if_int_compare:w \c_zero_int \c_zero_int \fi:
14392   =
14393   \__dim_compare_error:
14394 }

```

(End definition for `\dim_compare:nTF` and others. This function is documented on page 175.)

<pre> \dim_case:nn \dim_case:nnTF __dim_case:nnTF __dim_case:nw __dim_case_end:nw </pre>	<p>For dimension cases, the first task to fully expand the check condition. The over all idea is then much the same as for <code>\str_case:nn(TF)</code> as described in l3basics.</p> <pre> 14395 \cs_new:Npn \dim_case:nnTF #1 14396 { 14397 \exp:w 14398 \exp_args:Nf __dim_case:nnTF { \dim_eval:n {#1} } 14399 } 14400 \cs_new:Npn \dim_case:nnT #1#2#3 14401 { 14402 \exp:w 14403 \exp_args:Nf __dim_case:nnTF { \dim_eval:n {#1} } {#2} {#3} { } 14404 } 14405 \cs_new:Npn \dim_case:nnF #1#2 14406 { </pre>
---	---

```

14407     \exp:w
14408     \exp_args:Nf \__dim_case:nnTF { \dim_eval:n {#1} } {#2} { }
14409   }
14410   \cs_new:Npn \dim_case:nn #1#2
14411   {
14412     \exp:w
14413     \exp_args:Nf \__dim_case:nnTF { \dim_eval:n {#1} } {#2} { } { }
14414   }
14415   \cs_new:Npn \__dim_case:nnTF #1#2#3#4
14416   { \__dim_case:nw {#1} #2 {#1} { } \s__dim_mark {#3} \s__dim_mark {#4} \s__dim_stop }
14417   \cs_new:Npn \__dim_case:nw #1#2#3
14418   {
14419     \dim_compare:nNnTF {#1} = {#2}
14420     { \__dim_case_end:nw {#3} }
14421     { \__dim_case:nw {#1} }
14422   }
14423   \cs_new:Npn \__dim_case_end:nw #1#2#3 \s__dim_mark #4#5 \s__dim_stop
14424   { \exp_end: #1 #4 }

```

(End definition for `\dim_case:nnTF` and others. This function is documented on page 176.)

21.7 Dimension expression loops

`\dim_while_do:nn` `while_do` and `do_while` functions for dimensions. Same as for the `int` type only the names have changed.

```

\dim_until_do:nn
\dim_do_while:nn
\dim_do_until:nn
14425   \cs_new:Npn \dim_while_do:nn #1#2
14426   {
14427     \dim_compare:nT {#1}
14428     {
14429       #2
14430       \dim_while_do:nn {#1} {#2}
14431     }
14432   }
14433   \cs_new:Npn \dim_until_do:nn #1#2
14434   {
14435     \dim_compare:nF {#1}
14436     {
14437       #2
14438       \dim_until_do:nn {#1} {#2}
14439     }
14440   }
14441   \cs_new:Npn \dim_do_while:nn #1#2
14442   {
14443     #2
14444     \dim_compare:nT {#1}
14445     { \dim_do_while:nn {#1} {#2} }
14446   }
14447   \cs_new:Npn \dim_do_until:nn #1#2
14448   {
14449     #2
14450     \dim_compare:nF {#1}
14451     { \dim_do_until:nn {#1} {#2} }
14452   }

```

(End definition for `\dim_while_do:nN` and others. These functions are documented on page 177.)

`\dim_while_do:nNnn` `while_do` and `do_while` functions for dimensions. Same as for the `int` type only the names have changed.

```

\dim_until_do:nNnn
\dim_do_while:nNnn
\dim_do_until:nNnn
14453 \cs_new:Npn \dim_while_do:nNnn #1#2#3#4
14454 {
14455   \dim_compare:nNnT {#1} #2 {#3}
14456   {
14457     #4
14458     \dim_while_do:nNnn {#1} #2 {#3} {#4}
14459   }
14460 }
14461 \cs_new:Npn \dim_until_do:nNnn #1#2#3#4
14462 {
14463   \dim_compare:nNnF {#1} #2 {#3}
14464   {
14465     #4
14466     \dim_until_do:nNnn {#1} #2 {#3} {#4}
14467   }
14468 }
14469 \cs_new:Npn \dim_do_while:nNnn #1#2#3#4
14470 {
14471   #4
14472   \dim_compare:nNnT {#1} #2 {#3}
14473   { \dim_do_while:nNnn {#1} #2 {#3} {#4} }
14474 }
14475 \cs_new:Npn \dim_do_until:nNnn #1#2#3#4
14476 {
14477   #4
14478   \dim_compare:nNnF {#1} #2 {#3}
14479   { \dim_do_until:nNnn {#1} #2 {#3} {#4} }
14480 }

```

(End definition for `\dim_while_do:nNnn` and others. These functions are documented on page 177.)

21.8 Dimension step functions

`\dim_step_function:nnnN` Before all else, evaluate the initial value, step, and final value. Repeating a function by steps first needs a check on the direction of the steps. After that, do the function for the start value then step and loop around. It would be more symmetrical to test for a step size of zero before checking the sign, but we optimize for the most frequent case (positive step).

```

\__dim_step:wwwN
\__dim_step:NnnnN
14481 \cs_new:Npn \dim_step_function:nnnN #1#2#3
14482 {
14483   \exp_after:wN \__dim_step:wwwN
14484   \tex_the:D \__dim_eval:w #1 \exp_after:wN ;
14485   \tex_the:D \__dim_eval:w #2 \exp_after:wN ;
14486   \tex_the:D \__dim_eval:w #3 ;
14487 }
14488 \cs_new:Npn \__dim_step:wwwN #1; #2; #3; #4
14489 {
14490   \dim_compare:nNnTF {#2} > \c_zero_dim
14491   { \__dim_step:NnnnN > }

```

```

14492     {
14493         \dim_compare:nNnTF {#2} = \c_zero_dim
14494         {
14495             \__kernel_msg_expandable_error:nnn { kernel } { zero-step } {#4}
14496             \use_none:nnnn
14497         }
14498         { \__dim_step:NnnnN < }
14499     }
14500     {#1} {#2} {#3} #4
14501 }
14502 \cs_new:Npn \__dim_step:NnnnN #1#2#3#4#5
14503 {
14504     \dim_compare:nNf {#2} #1 {#4}
14505     {
14506         #5 {#2}
14507         \exp_args:NNf \__dim_step:NnnnN
14508         #1 { \dim_eval:n { #2 + #3 } } {#3} {#4} #5
14509     }
14510 }

```

(End definition for `\dim_step_function:nnnN`, `__dim_step:wwwN`, and `__dim_step:NnnnN`. This function is documented on page 177.)

```

\dim_step_inline:nnnn
\dim_step_variable:nnnNn
  \__dim_step:NNnnnn

```

The approach here is to build a function, with a global integer required to make the nesting safe (as seen in other in line functions), and map that function using `\dim_step_function:nnnN`. We put a `\prg_break_point:Nn` so that `map_break` functions from other modules correctly decrement `\g__kernel_prg_map_int` before looking for their own break point. The first argument is `\scan_stop:`, so that no breaking function recognizes this break point as its own.

```

14511 \cs_new_protected:Npn \dim_step_inline:nnnn
14512 {
14513     \int_gincr:N \g__kernel_prg_map_int
14514     \exp_args:NNc \__dim_step:NNnnnn
14515     \cs_gset_protected:Npn
14516     { \__dim_map_ \int_use:N \g__kernel_prg_map_int :w }
14517 }
14518 \cs_new_protected:Npn \dim_step_variable:nnnNn #1#2#3#4#5
14519 {
14520     \int_gincr:N \g__kernel_prg_map_int
14521     \exp_args:NNc \__dim_step:NNnnnn
14522     \cs_gset_protected:Npx
14523     { \__dim_map_ \int_use:N \g__kernel_prg_map_int :w }
14524     {#1}{#2}{#3}
14525     {
14526         \tl_set:Nn \exp_not:N #4 {##1}
14527         \exp_not:n {#5}
14528     }
14529 }
14530 \cs_new_protected:Npn \__dim_step:NNnnnn #1#2#3#4#5#6
14531 {
14532     #1 #2 ##1 {#6}
14533     \dim_step_function:nnnN {#3} {#4} {#5} #2
14534     \prg_break_point:Nn \scan_stop: { \int_gdecr:N \g__kernel_prg_map_int }
14535 }

```

(End definition for `\dim_step_inline:nnnn`, `\dim_step_variable:nnnNm`, and `_dim_step:NNnnnn`. These functions are documented on page 177.)

21.9 Using dim expressions and variables

`\dim_eval:n` Evaluating a dimension expression expandably.

```
14536 \cs_new:Npn \dim_eval:n #1
14537 { \dim_use:N \_dim_eval:w #1 \_dim_eval_end: }
```

(End definition for `\dim_eval:n`. This function is documented on page 178.)

`\dim_sign:n` See `\dim_abs:n`. Contrarily to `\int_sign:n` the case of a zero dimension cannot be distinguished from a positive dimension by looking only at the first character, since `0.2pt` and `0pt` start the same way. We need explicit comparisons. We start by distinguishing the most common case of a positive dimension.

```
14538 \cs_new:Npn \dim_sign:n #1
14539 {
14540   \int_value:w \exp_after:wN \_dim_sign:Nw
14541   \dim_use:N \_dim_eval:w #1 \_dim_eval_end: ;
14542   \exp_stop_f:
14543 }
14544 \cs_new:Npn \_dim_sign:Nw #1#2 ;
14545 {
14546   \if_dim:w #1#2 > \c_zero_dim
14547     1
14548   \else:
14549     \if_meaning:w - #1
14550       -1
14551     \else:
14552       0
14553     \fi:
14554   \fi:
14555 }
```

(End definition for `\dim_sign:n` and `_dim_sign:Nw`. This function is documented on page 178.)

`\dim_use:N` Accessing a $\langle dim \rangle$.

`\dim_use:c` `\cs_new_eq:NN \dim_use:N \tex_the:D`

We hand-code this for some speed gain:

```
14557 %\cs_generate_variant:Nn \dim_use:N { c }
14558 \cs_new:Npn \dim_use:c #1 { \tex_the:D \cs:w #1 \cs_end: }
```

(End definition for `\dim_use:N`. This function is documented on page 178.)

`\dim_to_decimal:n` A function which comes up often enough to deserve a place in the kernel. Evaluate the dimension expression `#1` then remove the trailing `pt`. When debugging is enabled, the argument is put in parentheses as this prevents the dimension expression from terminating early and leaving extra tokens lying around. This is used a lot by low-level manipulations.

```
14559 \cs_new:Npn \dim_to_decimal:n #1
14560 {
14561   \exp_after:wN
14562   \_dim_to_decimal:w \dim_use:N \_dim_eval:w #1 \_dim_eval_end:
14563 }
```

```

14564 \use:x
14565 {
14566   \cs_new:Npn \exp_not:N \_dim_to_decimal:w
14567     ##1 . ##2 \tl_to_str:n { pt }
14568 }
14569 {
14570   \int_compare:nNnTF {#2} > { 0 }
14571     { #1 . #2 }
14572     { #1 }
14573 }

```

(End definition for `\dim_to_decimal:n` and `_dim_to_decimal:w`. This function is documented on page 178.)

`\dim_to_decimal_in_bp:n` Conversion to big points is done using a scaling inside `_dim_eval:w` as ε -TEX does that using 64-bit precision. Here, 800/803 is the integer fraction for 72/72.27. This is a common case so is hand-coded for accuracy (and speed).

```

14574 \cs_new:Npn \dim_to_decimal_in_bp:n #1
14575 { \dim_to_decimal:n { ( #1 ) * 800 / 803 } }

```

(End definition for `\dim_to_decimal_in_bp:n`. This function is documented on page 179.)

`\dim_to_decimal_in_sp:n` Another hard-coded conversion: this one is necessary to avoid things going off-scale.

```

14576 \cs_new:Npn \dim_to_decimal_in_sp:n #1
14577 { \int_value:w \_dim_eval:w #1 \_dim_eval_end: }

```

(End definition for `\dim_to_decimal_in_sp:n`. This function is documented on page 179.)

`\dim_to_decimal_in_unit:nn` An analogue of `\dim_ratio:nn` that produces a decimal number as its result, rather than a rational fraction for use within dimension expressions.

```

14578 \cs_new:Npn \dim_to_decimal_in_unit:nn #1#2
14579 {
14580   \dim_to_decimal:n
14581   {
14582     1pt *
14583     \dim_ratio:nn {#1} {#2}
14584   }
14585 }

```

(End definition for `\dim_to_decimal_in_unit:nn`. This function is documented on page 179.)

`\dim_to_fp:n` Defined in `l3fp-convert`, documented here.

(End definition for `\dim_to_fp:n`. This function is documented on page 179.)

21.10 Viewing dim variables

`\dim_show:N` Diagnostics.

```

\dim_show:c 14586 \cs_new_eq:NN \dim_show:N \_kernel_register_show:N
14587 \cs_generate_variant:Nn \dim_show:N { c }

```

(End definition for `\dim_show:N`. This function is documented on page 179.)

\dim_show:n Diagnostics. We don't use the TeX primitive `\showthe` to show dimension expressions: this gives a more unified output.

```
14588 \cs_new_protected:Npn \dim_show:n
14589 { \msg_show_eval:Nn \dim_eval:n }
```

(End definition for `\dim_show:n`. This function is documented on page 180.)

\dim_log:N Diagnostics. Redirect output of `\dim_show:n` to the log.

```
\dim_log:c 14590 \cs_new_eq:NN \dim_log:N \__kernel_register_log:N
\dim_log:n 14591 \cs_new_eq:NN \dim_log:c \__kernel_register_log:c
14592 \cs_new_protected:Npn \dim_log:n
14593 { \msg_log_eval:Nn \dim_eval:n }
```

(End definition for `\dim_log:N` and `\dim_log:n`. These functions are documented on page 180.)

21.11 Constant dimensions

\c_zero_dim Constant dimensions.

```
\c_max_dim 14594 \dim_const:Nn \c_zero_dim { 0 pt }
14595 \dim_const:Nn \c_max_dim { 16383.99999 pt }
```

(End definition for `\c_zero_dim` and `\c_max_dim`. These variables are documented on page 180.)

21.12 Scratch dimensions

\l_tmpa_dim We provide two local and two global scratch registers, maybe we need more or less.

```
\l_tmpb_dim 14596 \dim_new:N \l_tmpa_dim
\g_tmpa_dim 14597 \dim_new:N \l_tmpb_dim
\g_tmpb_dim 14598 \dim_new:N \g_tmpa_dim
14599 \dim_new:N \g_tmpb_dim
```

(End definition for `\l_tmpa_dim` and others. These variables are documented on page 180.)

21.13 Creating and initialising skip variables

```
14600 <@@=skip>
```

\s_skip_stop Internal scan marks.

```
14601 \scan_new:N \s_skip_stop
```

(End definition for `\s_skip_stop`.)

\skip_new:N Allocation of a new internal registers.

```
\skip_new:c 14602 \cs_new_protected:Npn \skip_new:N #1
14603 {
14604   \__kernel_chk_if_free_cs:N #1
14605   \cs:w newskip \cs_end: #1
14606 }
14607 \cs_generate_variant:Nn \skip_new:N { c }
```

(End definition for `\skip_new:N`. This function is documented on page 180.)

\skip_const:Nn Contrarily to integer constants, we cannot avoid using a register, even for constants. See **\dim_const:Nn** for why we cannot use **\skip_gset:Nn**.

```

14608 \cs_new_protected:Npn \skip_const:Nn #1#2
14609 {
14610   \skip_new:N #1
14611   \tex_global:D #1 ~ \skip_eval:n {#2} \scan_stop:
14612 }
14613 \cs_generate_variant:Nn \skip_const:Nn { c }

```

(End definition for **\skip_const:Nn**. This function is documented on page 181.)

\skip_zero:N Reset the register to zero.

```

\skip_zero:c 14614 \cs_new_protected:Npn \skip_zero:N #1 { #1 \c_zero_skip }
\skip_gzero:N 14615 \cs_new_protected:Npn \skip_gzero:N #1 { \tex_global:D #1 \c_zero_skip }
\skip_gzero:c 14616 \cs_generate_variant:Nn \skip_zero:N { c }
               14617 \cs_generate_variant:Nn \skip_gzero:N { c }

```

(End definition for **\skip_zero:N** and **\skip_gzero:N**. These functions are documented on page 181.)

\skip_zero_new:N Create a register if needed, otherwise clear it.

```

\skip_zero_new:c 14618 \cs_new_protected:Npn \skip_zero_new:N #1
\skip_gzero_new:N 14619 { \skip_if_exist:NTF #1 { \skip_zero:N #1 } { \skip_new:N #1 } }
\skip_gzero_new:c 14620 \cs_new_protected:Npn \skip_gzero_new:N #1
                  14621 { \skip_if_exist:NTF #1 { \skip_gzero:N #1 } { \skip_new:N #1 } }
                  14622 \cs_generate_variant:Nn \skip_zero_new:N { c }
                  14623 \cs_generate_variant:Nn \skip_gzero_new:N { c }

```

(End definition for **\skip_zero_new:N** and **\skip_gzero_new:N**. These functions are documented on page 181.)

\skip_if_exist_p:N Copies of the **cs** functions defined in **l3basics**.

```

\skip_if_exist_p:c 14624 \prg_new_eq_conditional:NNn \skip_if_exist:N \cs_if_exist:N
\skip_if_exist:N\TF 14625 { TF , T , F , p }
\skip_if_exist:c\TF 14626 \prg_new_eq_conditional:NNn \skip_if_exist:c \cs_if_exist:c
                  14627 { TF , T , F , p }

```

(End definition for **\skip_if_exist:N\TF**. This function is documented on page 181.)

21.14 Setting skip variables

\skip_set:Nn Much the same as for dimensions.

```

\skip_set:cn 14628 \cs_new_protected:Npn \skip_set:Nn #1#2
\skip_gset:Nn 14629 { #1 ~ \tex_glueexpr:D #2 \scan_stop: }
\skip_gset:cn 14630 \cs_new_protected:Npn \skip_gset:Nn #1#2
              14631 { \tex_global:D #1 ~ \tex_glueexpr:D #2 \scan_stop: }
              14632 \cs_generate_variant:Nn \skip_set:Nn { c }
              14633 \cs_generate_variant:Nn \skip_gset:Nn { c }

```

(End definition for **\skip_set:Nn** and **\skip_gset:Nn**. These functions are documented on page 181.)

\skip_set_eq:NN All straightforward.

```

\skip_set_eq:cn 14634 \cs_new_protected:Npn \skip_set_eq:NN #1#2 { #1 = #2 }
\skip_set_eq:Nc 14635 \cs_generate_variant:Nn \skip_set_eq:NN { c , Nc , cc }
\skip_set_eq:cc 14636 \cs_new_protected:Npn \skip_gset_eq:NN #1#2 { \tex_global:D #1 = #2 }
\skip_gset_eq:NN 14637 \cs_generate_variant:Nn \skip_gset_eq:NN { c , Nc , cc }
\skip_gset_eq:cn
\skip_gset_eq:Nc
\skip_gset_eq:cc

```

(End definition for `\skip_set_eq:Nn` and `\skip_gset_eq:Nn`. These functions are documented on page 181.)

```

\skip_add:Nn Using by here deals with the (incorrect) case \skip123.
\skip_add:cn 14638 \cs_new_protected:Npn \skip_add:Nn #1#2
\skip_gadd:Nn 14639 { \tex_advance:D #1 by \tex_glueexpr:D #2 \scan_stop: }
\skip_gadd:cn 14640 \cs_new_protected:Npn \skip_gadd:Nn #1#2
\skip_sub:Nn 14641 { \tex_global:D \tex_advance:D #1 by \tex_glueexpr:D #2 \scan_stop: }
\skip_sub:cn 14642 \cs_generate_variant:Nn \skip_add:Nn { c }
\skip_gsub:Nn 14643 \cs_generate_variant:Nn \skip_gadd:Nn { c }
\skip_gsub:cn 14644 \cs_new_protected:Npn \skip_sub:Nn #1#2
14645 { \tex_advance:D #1 by - \tex_glueexpr:D #2 \scan_stop: }
14646 \cs_new_protected:Npn \skip_gsub:Nn #1#2
14647 { \tex_global:D \tex_advance:D #1 by - \tex_glueexpr:D #2 \scan_stop: }
14648 \cs_generate_variant:Nn \skip_sub:Nn { c }
14649 \cs_generate_variant:Nn \skip_gsub:Nn { c }

```

(End definition for `\skip_add:Nn` and others. These functions are documented on page 181.)

21.15 Skip expression conditionals

`\skip_if_eq_p:n` Comparing skips means doing two expansions to make strings, and then testing them.
`\skip_if_eq:nnTF` As a result, only equality is tested.

```

14650 \prg_new_conditional:Npnn \skip_if_eq:nn #1#2 { p , T , F , TF }
14651 {
14652   \str_if_eq:eeTF { \skip_eval:n { #1 } } { \skip_eval:n { #2 } }
14653   { \prg_return_true: }
14654   { \prg_return_false: }
14655 }

```

(End definition for `\skip_if_eq:nnTF`. This function is documented on page 182.)

`\skip_if_finite_p:n` With ε -TeX, we have an easy access to the order of infinities of the stretch and shrink components of a skip. However, to access both, we either need to evaluate the expression twice, or evaluate it, then call an auxiliary to extract both pieces of information from the result. Since we are going to need an auxiliary anyways, it is quicker to make it search for the string `fil` which characterizes infinite glue.
`\skip_if_finite:nTF`
`__skip_if_finite:wwNw`

```

14656 \cs_set_protected:Npn \__skip_tmp:w #1
14657 {
14658   \prg_new_conditional:Npnn \skip_if_finite:n ##1 { p , T , F , TF }
14659   {
14660     \exp_after:wN \__skip_if_finite:wwNw
14661     \skip_use:N \tex_glueexpr:D ##1 ; \prg_return_false:
14662     #1 ; \prg_return_true: \s__skip_stop
14663   }
14664   \cs_new:Npn \__skip_if_finite:wwNw ##1 #1 ##2 ; ##3 ##4 \s__skip_stop {##3}
14665   }
14666   \exp_args:No \__skip_tmp:w { \tl_to_str:n { fil } }

```

(End definition for `\skip_if_finite:nTF` and `__skip_if_finite:wwNw`. This function is documented on page 182.)

21.16 Using skip expressions and variables

`\skip_eval:n` Evaluating a skip expression expandably.

```
14667 \cs_new:Npn \skip_eval:n #1
14668 { \skip_use:N \tex_glueexpr:D #1 \scan_stop: }
```

(End definition for `\skip_eval:n`. This function is documented on page 182.)

`\skip_use:N` Accessing a `\skip`.

```
\skip_use:c
14669 \cs_new_eq:NN \skip_use:N \tex_the:D
14670 %\cs_generate_variant:Nn \skip_use:N { c }
14671 \cs_new:Npn \skip_use:c #1 { \tex_the:D \cs:w #1 \cs_end: }
```

(End definition for `\skip_use:N`. This function is documented on page 182.)

21.17 Inserting skips into the output

`\skip_horizontal:N` Inserting skips.

```
\skip_horizontal:c
\skip_horizontal:n
\skip_vertical:N
\skip_vertical:c
\skip_vertical:n
14672 \cs_new_eq:NN \skip_horizontal:N \tex_hskip:D
14673 \cs_new:Npn \skip_horizontal:n #1
14674 { \skip_horizontal:N \tex_glueexpr:D #1 \scan_stop: }
14675 \cs_new_eq:NN \skip_vertical:N \tex_vskip:D
14676 \cs_new:Npn \skip_vertical:n #1
14677 { \skip_vertical:N \tex_glueexpr:D #1 \scan_stop: }
14678 \cs_generate_variant:Nn \skip_horizontal:N { c }
14679 \cs_generate_variant:Nn \skip_vertical:N { c }
```

(End definition for `\skip_horizontal:N` and others. These functions are documented on page 183.)

21.18 Viewing skip variables

`\skip_show:N` Diagnostics.

```
\skip_show:c
14680 \cs_new_eq:NN \skip_show:N \__kernel_register_show:N
14681 \cs_generate_variant:Nn \skip_show:N { c }
```

(End definition for `\skip_show:N`. This function is documented on page 182.)

`\skip_show:n` Diagnostics. We don't use the TeX primitive `\showthe` to show skip expressions: this gives a more unified output.

```
14682 \cs_new_protected:Npn \skip_show:n
14683 { \msg_show_eval:Nn \skip_eval:n }
```

(End definition for `\skip_show:n`. This function is documented on page 182.)

`\skip_log:N` Diagnostics. Redirect output of `\skip_show:n` to the log.

```
\skip_log:c
\skip_log:n
14684 \cs_new_eq:NN \skip_log:N \__kernel_register_log:N
14685 \cs_new_eq:NN \skip_log:c \__kernel_register_log:c
14686 \cs_new_protected:Npn \skip_log:n
14687 { \msg_log_eval:Nn \skip_eval:n }
```

(End definition for `\skip_log:N` and `\skip_log:n`. These functions are documented on page 183.)

21.19 Constant skips

`\c_zero_skip` Skips with no rubber component are just dimensions but need to terminate correctly.

```
\c_max_skip 14688 \skip_const:Nn \c_zero_skip { \c_zero_dim }
14689 \skip_const:Nn \c_max_skip { \c_max_dim }
```

(End definition for `\c_zero_skip` and `\c_max_skip`. These functions are documented on page 183.)

21.20 Scratch skips

`\l_tmpa_skip` We provide two local and two global scratch registers, maybe we need more or less.

```
\l_tmpb_skip 14690 \skip_new:N \l_tmpa_skip
\g_tmpa_skip 14691 \skip_new:N \l_tmpb_skip
\g_tmpb_skip 14692 \skip_new:N \g_tmpa_skip
14693 \skip_new:N \g_tmpb_skip
```

(End definition for `\l_tmpa_skip` and others. These variables are documented on page 183.)

21.21 Creating and initialising muskip variables

`\muskip_new:N` And then we add muskips.

```
\muskip_new:c 14694 \cs_new_protected:Npn \muskip_new:N #1
14695 {
14696   \__kernel_chk_if_free_cs:N #1
14697   \cs:w newmuskip \cs_end: #1
14698 }
14699 \cs_generate_variant:Nn \muskip_new:N { c }
```

(End definition for `\muskip_new:N`. This function is documented on page 184.)

`\muskip_const:Nn` See `\skip_const:Nn`.

```
\muskip_const:cn 14700 \cs_new_protected:Npn \muskip_const:Nn #1#2
14701 {
14702   \muskip_new:N #1
14703   \tex_global:D #1 ~ \muskip_eval:n {#2} \scan_stop:
14704 }
14705 \cs_generate_variant:Nn \muskip_const:Nn { c }
```

(End definition for `\muskip_const:Nn`. This function is documented on page 184.)

`\muskip_zero:N` Reset the register to zero.

```
\muskip_zero:c 14706 \cs_new_protected:Npn \muskip_zero:N #1
\muskip_gzero:N 14707 { #1 \c_zero_muskip }
\muskip_gzero:c 14708 \cs_new_protected:Npn \muskip_gzero:N #1
14709 { \tex_global:D #1 \c_zero_muskip }
14710 \cs_generate_variant:Nn \muskip_zero:N { c }
14711 \cs_generate_variant:Nn \muskip_gzero:N { c }
```

(End definition for `\muskip_zero:N` and `\muskip_gzero:N`. These functions are documented on page 184.)

`\muskip_zero_new:N` Create a register if needed, otherwise clear it.

```

\muskip_zero_new:c 14712 \cs_new_protected:Npn \muskip_zero_new:N #1
\muskip_gzero_new:N 14713 { \muskip_if_exist:NTF #1 { \muskip_zero:N #1 } { \muskip_new:N #1 } }
\muskip_gzero_new:c 14714 \cs_new_protected:Npn \muskip_gzero_new:N #1
14715 { \muskip_if_exist:NTF #1 { \muskip_gzero:N #1 } { \muskip_new:N #1 } }
14716 \cs_generate_variant:Nn \muskip_zero_new:N { c }
14717 \cs_generate_variant:Nn \muskip_gzero_new:N { c }

(End definition for \muskip_zero_new:N and \muskip_gzero_new:N. These functions are documented on
page 184.)

```

`\muskip_if_exist_p:N` Copies of the cs functions defined in l3basics.

```

\muskip_if_exist_p:c 14718 \prg_new_eq_conditional:NNn \muskip_if_exist:N \cs_if_exist:N
\muskip_if_exist:NTF 14719 { TF , T , F , p }
\muskip_if_exist:cTF 14720 \prg_new_eq_conditional:NNn \muskip_if_exist:c \cs_if_exist:c
14721 { TF , T , F , p }

(End definition for \muskip_if_exist:NTF. This function is documented on page 184.)

```

21.22 Setting muskip variables

`\muskip_set:Nn` This should be pretty familiar.

```

\muskip_set:cn 14722 \cs_new_protected:Npn \muskip_set:Nn #1#2
\muskip_gset:Nn 14723 { #1 ~ \tex_muexpr:D #2 \scan_stop: }
\muskip_gset:cn 14724 \cs_new_protected:Npn \muskip_gset:Nn #1#2
14725 { \tex_global:D #1 ~ \tex_muexpr:D #2 \scan_stop: }
14726 \cs_generate_variant:Nn \muskip_set:Nn { c }
14727 \cs_generate_variant:Nn \muskip_gset:Nn { c }

(End definition for \muskip_set:Nn and \muskip_gset:Nn. These functions are documented on page
185.)

```

`\muskip_set_eq:NN` All straightforward.

```

\muskip_set_eq:cN 14728 \cs_new_protected:Npn \muskip_set_eq:NN #1#2 { #1 = #2 }
\muskip_set_eq:Nc 14729 \cs_generate_variant:Nn \muskip_set_eq:NN { c , Nc , cc }
\muskip_set_eq:cc 14730 \cs_new_protected:Npn \muskip_gset_eq:NN #1#2 { \tex_global:D #1 = #2 }
\muskip_gset_eq:NN 14731 \cs_generate_variant:Nn \muskip_gset_eq:NN { c , Nc , cc }
\muskip_gset_eq:cN
\muskip_gset_eq:Nc
\muskip_gset_eq:cc

(End definition for \muskip_set_eq:NN and \muskip_gset_eq:NN. These functions are documented on
page 185.)

```

`\muskip_add:Nn` Using by here deals with the (incorrect) case `\muskip123`.

```

\muskip_add:cn 14732 \cs_new_protected:Npn \muskip_add:Nn #1#2
\muskip_gadd:Nn 14733 { \tex_advance:D #1 by \tex_muexpr:D #2 \scan_stop: }
\muskip_gadd:cn 14734 \cs_new_protected:Npn \muskip_gadd:Nn #1#2
\muskip_sub:Nn 14735 { \tex_global:D \tex_advance:D #1 by \tex_muexpr:D #2 \scan_stop: }
\muskip_sub:cn 14736 \cs_generate_variant:Nn \muskip_add:Nn { c }
\muskip_gsub:Nn 14737 \cs_generate_variant:Nn \muskip_gadd:Nn { c }
\muskip_gsub:cn 14738 \cs_new_protected:Npn \muskip_sub:Nn #1#2
14739 { \tex_advance:D #1 by - \tex_muexpr:D #2 \scan_stop: }
14740 \cs_new_protected:Npn \muskip_gsub:Nn #1#2
14741 { \tex_global:D \tex_advance:D #1 by - \tex_muexpr:D #2 \scan_stop: }
14742 \cs_generate_variant:Nn \muskip_sub:Nn { c }
14743 \cs_generate_variant:Nn \muskip_gsub:Nn { c }

(End definition for \muskip_add:Nn and others. These functions are documented on page 184.)

```

21.23 Using muskip expressions and variables

\muskip_eval:n Evaluating a muskip expression expandably.

```
14744 \cs_new:Npn \muskip_eval:n #1
14745 { \muskip_use:N \tex_muexpr:D #1 \scan_stop: }
```

(End definition for \muskip_eval:n. This function is documented on page 185.)

\muskip_use:N Accessing a $\langle muskip \rangle$.
\muskip_use:c

```
14746 \cs_new_eq:NN \muskip_use:N \tex_the:D
14747 \cs_generate_variant:Nn \muskip_use:N { c }
```

(End definition for \muskip_use:N. This function is documented on page 185.)

21.24 Viewing muskip variables

\muskip_show:N Diagnostics.
\muskip_show:c

```
14748 \cs_new_eq:NN \muskip_show:N \__kernel_register_show:N
14749 \cs_generate_variant:Nn \muskip_show:N { c }
```

(End definition for \muskip_show:N. This function is documented on page 185.)

\muskip_show:n Diagnostics. We don't use the TeX primitive \showthe to show muskip expressions: this gives a more unified output.

```
14750 \cs_new_protected:Npn \muskip_show:n
14751 { \msg_show_eval:Nn \muskip_eval:n }
```

(End definition for \muskip_show:n. This function is documented on page 186.)

\muskip_log:N Diagnostics. Redirect output of \muskip_show:n to the log.
\muskip_log:c
\muskip_log:n

```
14752 \cs_new_eq:NN \muskip_log:N \__kernel_register_log:N
14753 \cs_new_eq:NN \muskip_log:c \__kernel_register_log:c
14754 \cs_new_protected:Npn \muskip_log:n
14755 { \msg_log_eval:Nn \muskip_eval:n }
```

(End definition for \muskip_log:N and \muskip_log:n. These functions are documented on page 186.)

21.25 Constant muskips

\c_zero_muskip Constant muskips given by their value.
\c_max_muskip

```
14756 \muskip_const:Nn \c_zero_muskip { 0 mu }
14757 \muskip_const:Nn \c_max_muskip { 16383.99999 mu }
```

(End definition for \c_zero_muskip and \c_max_muskip. These functions are documented on page 186.)

21.26 Scratch muskips

\l_tmpa_muskip We provide two local and two global scratch registers, maybe we need more or less.
\l_tmpb_muskip
\g_tmpa_muskip
\g_tmpb_muskip

```
14758 \muskip_new:N \l_tmpa_muskip
14759 \muskip_new:N \l_tmpb_muskip
14760 \muskip_new:N \g_tmpa_muskip
14761 \muskip_new:N \g_tmpb_muskip
```

(End definition for \l_tmpa_muskip and others. These variables are documented on page 186.)

```
14762 \endpackage
```

22 l3keys Implementation

14763 $\langle *package \rangle$

22.1 Low-level interface

The low-level key parser’s implementation is based heavily on `expkv`. Compared to `keyval` it adds a number of additional “safety” requirements and allows to process the parsed list of key–value pairs in a variety of ways. The net result is that this code needs around one and a half the amount of time as `keyval` to parse the same list of keys. To optimise speed as far as reasonably practical, a number of lower-level approaches are taken rather than using the higher-level `expl3` interfaces.

14764 $\langle @@=keyval \rangle$

```
\s__keyval_nil
\s__keyval_mark
\s__keyval_stop
\s__keyval_tail
14765 \scan_new:N \s__keyval_nil
14766 \scan_new:N \s__keyval_mark
14767 \scan_new:N \s__keyval_stop
14768 \scan_new:N \s__keyval_tail
```

(End definition for `\s__keyval_nil` and others.)

This temporary macro will be used since some of the definitions will need an active comma or equals sign. Inside of this macro `#1` will be the active comma and `#2` will be the active equals sign.

```
14769 \group_begin:
14770   \cs_set_protected:Npn \__keyval_tmp:NN #1#2
14771   {
```

`\keyval_parse:nnn`
`\keyval_parse:NNn`

The main function starts the first of two loops. The outer loop splits the key–value list at active commas, the inner loop will do so at other commas. The use of `\s__keyval_mark` here prevents loss of braces from the key argument.

```
14772   \cs_new:Npn \keyval_parse:nnn ##1 ##2 ##3
14773   { \__keyval_loop_active:nnw {##1} {##2} \s__keyval_mark ##3 #1 \s__keyval_tail #1 }
14774   \cs_new_eq:NN \keyval_parse:NNn \keyval_parse:nnn
```

(End definition for `\keyval_parse:nnn` and `\keyval_parse:NNn`. These functions are documented on page 199.)

`__keyval_loop_active:nnw`

First a fast test for the end of the loop is done, it’ll gobble everything up to a `\s__keyval_tail`. The loop ending macro will gobble everything to the last comma in this definition. If the end isn’t reached yet, start the second loop splitting at other commas, the next iteration of this first loop will be inserted by the end of `__keyval_loop_active:nnw`.

```
14775   \cs_new:Npn \__keyval_loop_active:nnw ##1 ##2 ##3 #1
14776   {
14777     \__keyval_if_recursion_tail:w ##3
14778     \__keyval_end_loop_active:w \s__keyval_tail
14779     \__keyval_loop_other:nnw {##1} {##2} ##3 , \s__keyval_tail ,
14780   }
```

(End definition for `__keyval_loop_active:nnw`.)

`__keyval_split_other:w` These two macros allow to split at the first equals sign of category 12 or 13. At the same time they also execute branching by inserting the first token following `\s__keyval_mark` that followed the equals sign. Hence they also test for the presence of such an equals sign simultaneously.

```

14781 \cs_new:Npn \__keyval_split_other:w ##1 = ##2 \s__keyval_mark ##3
14782 { ##3 ##1 \s__keyval_stop \s__keyval_mark ##2 }
14783 \cs_new:Npn \__keyval_split_active:w ##1 #2 ##2 \s__keyval_mark ##3
14784 { ##3 ##1 \s__keyval_stop \s__keyval_mark ##2 }

```

(End definition for `__keyval_split_other:w` and `__keyval_split_active:w`.)

`__keyval_loop_other:nnw` The second loop uses the same test for its end as the first loop, next it splits at the first active equals sign using `__keyval_split_active:w`. The `\s__keyval_nil` prevents accidental brace stripping and acts as a delimiter in the next steps. First testing for an active equals sign will reduce the number of necessary expansion steps for the expected average use case of other equals signs and hence perform better on average.

```

14785 \cs_new:Npn \__keyval_loop_other:nnw ##1 ##2 ##3 ,
14786 {
14787   \__keyval_if_recursion_tail:w ##3
14788   \__keyval_end_loop_other:w \s__keyval_tail
14789   \__keyval_split_active:w ##3 \s__keyval_nil
14790   \s__keyval_mark \__keyval_split_active_auxi:w
14791   #2 \s__keyval_mark \__keyval_clean_up_active:w
14792   {##1} {##2}
14793   \s__keyval_mark
14794 }

```

(End definition for `__keyval_loop_other:nnw`.)

`__keyval_split_active_auxi:w` After `__keyval_split_active:w` the following will only be called if there was at least one active equals sign in the current key–value pair. Therefore this is the execution branch for a key–value pair with an active equals sign. `##1` will be everything up to the first active equals sign. First it tests for other equals signs in the key name, which will eventually throw an error via `__keyval_misplaced_equal_after_active_error:w`. If none was found we forward the key to `__keyval_split_active_auxii:w`.

```

14795 \cs_new:Npn \__keyval_split_active_auxi:w ##1 \s__keyval_stop
14796 {
14797   \__keyval_split_other:w ##1 \s__keyval_nil
14798   \s__keyval_mark \__keyval_misplaced_equal_after_active_error:w
14799   = \s__keyval_mark \__keyval_split_active_auxii:w
14800 }

```

`__keyval_split_active_auxii:w` gets the correct key name with a leading `\s__keyval_mark` as `##1`. It has to sanitise the remainder of the previous test and trims the key name which will be forwarded to `__keyval_split_active_auxiii:w`.

```

14801 \cs_new:Npn \__keyval_split_active_auxii:w
14802   ##1 \s__keyval_nil \s__keyval_mark \__keyval_misplaced_equal_after_active_error:w
14803   \s__keyval_stop \s__keyval_mark
14804   ##2 \s__keyval_nil #2 \s__keyval_mark \__keyval_clean_up_active:w
14805   { \__keyval_trim:nN {##1} \__keyval_split_active_auxiii:w ##2 \s__keyval_nil }

```


Next we test for a misplaced active equals sign in the value, if none is found `__keyval_split_active_auxiv:w` will be called.

```

14806     \cs_new:Npn \__keyval_split_active_auxiii:w ##1 ##2 \s__keyval_nil
14807     {
14808         \__keyval_split_active:w ##2 \s__keyval_nil
14809         \s__keyval_mark \__keyval_misplaced_equal_in_split_error:w
14810         #2 \s__keyval_mark \__keyval_split_active_auxiv:w
14811         {##1}
14812     }

```

This runs the last test after sanitising the remainder of the previous one. This time test for a misplaced equals sign of category 12 in the value. Finally the last auxiliary macro will be called.

```

14813     \cs_new:Npn \__keyval_split_active_auxiv:w
14814         ##1 \s__keyval_nil \s__keyval_mark \__keyval_misplaced_equal_in_split_error:w
14815         \s__keyval_stop \s__keyval_mark
14816     {
14817         \__keyval_split_other:w ##1 \s__keyval_nil
14818         \s__keyval_mark \__keyval_misplaced_equal_in_split_error:w
14819         = \s__keyval_mark \__keyval_split_active_auxv:w
14820     }

```

This last macro in this execution branch sanitises the last test, trims the value and passes it to `__keyval_pair:nnnn`.

```

14821     \cs_new:Npn \__keyval_split_active_auxv:w
14822         ##1 \s__keyval_nil \s__keyval_mark \__keyval_misplaced_equal_in_split_error:w
14823         \s__keyval_stop \s__keyval_mark
14824         { \__keyval_trim:nN { ##1 } \__keyval_pair:nnnn }

```

(End definition for __keyval_split_active_auxi:w and others.)

`__keyval_clean_up_active:w` The following is the branch taken if the key–value pair doesn’t contain an active equals sign. The remainder of that test will be cleaned up by `__keyval_clean_up_active:w` which will then split at an equals sign of category other.

```

14825     \cs_new:Npn \__keyval_clean_up_active:w
14826         ##1 \s__keyval_nil \s__keyval_mark \__keyval_split_active_auxi:w \s__keyval_stop
14827     {
14828         \__keyval_split_other:w ##1 \s__keyval_nil
14829         \s__keyval_mark \__keyval_split_other_auxi:w
14830         = \s__keyval_mark \__keyval_clean_up_other:w
14831     }

```

(End definition for __keyval_clean_up_active:w.)

`__keyval_split_other_auxi:w` This is executed if the key–value pair doesn’t contain an active equals sign but at least one other. `##1` of `__keyval_split_other_auxi:w` will contain the complete key name, which is trimmed and forwarded to the next auxiliary macro.

`__keyval_split_other_auxii:w`
`__keyval_split_other_auxiii:w`

```

14832     \cs_new:Npn \__keyval_split_other_auxi:w ##1 \s__keyval_stop
14833     { \__keyval_trim:nN { ##1 } \__keyval_split_other_auxii:w }

```

We know that the value doesn’t contain misplaced active equals signs but we have to test for others. Also we need to sanitise the previous test, which is done here and not earlier to avoid superfluous argument grabbing.

```

14834     \cs_new:Npn \__keyval_split_other_auxii:w

```

```

14835         ##1 ##2 \s__keyval_nil = \s__keyval_mark \_keyval_clean_up_other:w
14836     {
14837         \_keyval_split_other:w ##2 \s__keyval_nil
14838         \s__keyval_mark \_keyval_misplaced_equal_in_split_error:w
14839         = \s__keyval_mark \_keyval_split_other_auxiii:w
14840         { ##1 }
14841     }

```

_keyval_split_other_auxiii:w sanitises the test for other equals signs, trims the value and forwards it to _keyval_pair:nnnn.

```

14842     \cs_new:Npn \_keyval_split_other_auxiii:w
14843         ##1 \s__keyval_nil \s__keyval_mark \_keyval_misplaced_equal_in_split_error:w
14844         \s__keyval_stop \s__keyval_mark
14845     { \_keyval_trim:nN { ##1 } \_keyval_pair:nnnn }

```

(End definition for _keyval_split_other_auxi:w, _keyval_split_other_auxii:w, and _keyval_split_other_auxiii:w.)

_keyval_clean_up_other:w _keyval_clean_up_other:w is the last branch that might exist. It is called if no equals sign was found, hence the only possibilities left are a blank list element, which is to be skipped, or a lonely key. If it's no empty list element this will trim the key name and forward it to _keyval_key:nn.

```

14846     \cs_new:Npn \_keyval_clean_up_other:w
14847         ##1 \s__keyval_nil \s__keyval_mark \_keyval_split_other_auxi:w \s__keyval_stop \
14848     {
14849         \_keyval_if_blank:w ##1 \s__keyval_nil \s__keyval_stop \_keyval_blank_true:w
14850         \s__keyval_mark \s__keyval_stop
14851         \_keyval_trim:nN { ##1 } \_keyval_key:nn
14852     }

```

(End definition for _keyval_clean_up_other:w.)

keyval_misplaced_equal_after_active_error:w All these two macros do is gobble the remainder of the current other loop execution and throw an error. Afterwards they have to insert the next loop iteration.

_keyval_misplaced_equal_in_split_error:w

```

14853     \cs_new:Npn \_keyval_misplaced_equal_after_active_error:w
14854         \s__keyval_mark ##1 \s__keyval_stop \s__keyval_mark ##2 \s__keyval_nil
14855         = \s__keyval_mark \_keyval_split_active_auxii:w
14856         \s__keyval_mark ##3 \s__keyval_nil
14857         #2 \s__keyval_mark \_keyval_clean_up_active:w
14858     {
14859         \_kernel_msg_expandable_error:nn
14860         { kernel } { misplaced-equals-sign }
14861         \_keyval_loop_other:nnw
14862     }
14863     \cs_new:Npn \_keyval_misplaced_equal_in_split_error:w
14864         \s__keyval_mark ##1 \s__keyval_stop \s__keyval_mark ##2 \s__keyval_nil
14865         ##3 \s__keyval_mark ##4 ##5
14866     {
14867         \_kernel_msg_expandable_error:nn
14868         { kernel } { misplaced-equals-sign }
14869         \_keyval_loop_other:nnw
14870     }

```

(End definition for _keyval_misplaced_equal_after_active_error:w and _keyval_misplaced_equal_in_split_error:w.)

`__keyval_end_loop_other:w` All that's left for the parsing loops are the macros which end the recursion. Both just gobble the remaining tokens of the respective loop including the next recursion call. `__keyval_end_loop_active:w` `__keyval_end_loop_other:w` also has to insert the next iteration of the active loop.

```

14871 \cs_new:Npn \__keyval_end_loop_other:w
14872 \s__keyval_tail
14873 \__keyval_split_active:w
14874 \s__keyval_mark \s__keyval_tail
14875 \s__keyval_nil \s__keyval_mark
14876 \__keyval_split_active_auxi:w
14877 #2 \s__keyval_mark \__keyval_clean_up_active:w
14878 { \__keyval_loop_active:nnw }
14879 \cs_new:Npn \__keyval_end_loop_active:w
14880 \s__keyval_tail
14881 \__keyval_loop_other:nnw ##1 \s__keyval_mark \s__keyval_tail , \s__keyval_tail ,
14882 { }

```

(End definition for `__keyval_end_loop_other:w` and `__keyval_end_loop_active:w`.)

The parsing loops are done, so here ends the definition of `__keyval_tmp:NN`, which will finally set up the macros.

```

14883 }
14884 \char_set_catcode_active:n { '\ , }
14885 \char_set_catcode_active:n { '\ = }
14886 \__keyval_tmp:NN , =
14887 \group_end:

```

`__keyval_pair:nnnn` These macros will be called on the parsed keys and values of the key–value list. All arguments are completely trimmed. They test for blank key names and call the functions passed to `\keyval_parse:nnn` inside of `\exp_not:n` with the correct arguments. Afterwards they insert the next iteration of the other loop.

```

14888 \cs_new:Npn \__keyval_pair:nnnn #1 #2 #3 #4
14889 {
14890 \__keyval_if_blank:w \s__keyval_mark #2 \s__keyval_nil \s__keyval_stop \__keyval_blank_
14891 \s__keyval_mark \s__keyval_stop
14892 \exp_not:n { #4 { #2 } { #1 } }
14893 \__keyval_loop_other:nnw {#3} {#4}
14894 }
14895 \cs_new:Npn \__keyval_key:nn #1 #2
14896 {
14897 \__keyval_if_blank:w \s__keyval_mark #1 \s__keyval_nil \s__keyval_stop \__keyval_blank_
14898 \s__keyval_mark \s__keyval_stop
14899 \exp_not:n { #2 { #1 } }
14900 \__keyval_loop_other:nnw {#2}
14901 }

```

(End definition for `__keyval_pair:nnnn` and `__keyval_key:nn`.)

`__keyval_if_empty:w` All these tests work by gobbling tokens until a certain combination is met, which makes them pretty fast. The test for a blank argument should be called with an arbitrary token following the argument. Each of these utilize the fact that the argument will contain a leading `\s__keyval_mark`.

```

14902 \cs_new:Npn \__keyval_if_empty:w #1 \s__keyval_mark \s__keyval_stop { }
14903 \cs_new:Npn \__keyval_if_blank:w \s__keyval_mark #1 { \__keyval_if_empty:w \s__keyval_mark
14904 \cs_new:Npn \__keyval_if_recursion_tail:w \s__keyval_mark #1 \s__keyval_tail { }

```

(End definition for `_keyval_if_empty:w`, `_keyval_if_blank:w`, and `_keyval_if_recursion_tail:w`.)

`_keyval_blank_true:w`
`_keyval_blank_key_error:w`

These macros will be called if the tests above didn't gobble them, they execute the branching.

```
14905 \cs_new:Npn \_keyval_blank_true:w \s_keyval_mark \s_keyval_stop \_keyval_trim:nN #1 \_
14906 { \_keyval_loop_other:nnw }
14907 \cs_new:Npn \_keyval_blank_key_error:w \s_keyval_mark \s_keyval_stop \exp_not:n #1
14908 {
14909   \_kernel_msg_expandable_error:nn
14910   { kernel } { blank-key-name }
14911 }
```

(End definition for `_keyval_blank_true:w` and `_keyval_blank_key_error:w`.)

Two messages for the low level parsing system.

```
14912 \_kernel_msg_new:nnn { kernel } { misplaced-equals-sign }
14913 { Misplaced~equals~sign~in~key~value~input~\msg_line_context: }
14914 \_kernel_msg_new:nnn { kernel } { blank-key-name }
14915 { Blank~key~name~in~key~value~input~\msg_line_context: }
```

`_keyval_trim:nN`
`_keyval_trim_auxi:w`
`_keyval_trim_auxii:w`
`_keyval_trim_auxiii:w`
`_keyval_trim_auxiv:w`

And an adapted version of `_tl_trim_spaces:nn` which is a bit faster for our use case, as it can strip the braces at the end. This is pretty much the same concept, so I won't comment on it here. The speed gain by using this instead of `\tl_trim_spaces_apply:nN` is about 10 % of the total time for `\keyval_parse:NNn` with one key and one key-value pair, so I think it's worth it.

```
14916 \group_begin:
14917 \cs_set_protected:Npn \_keyval_tmp:n #1
14918 {
14919   \cs_new:Npn \_keyval_trim:nN ##1
14920   {
14921     \_keyval_trim_auxi:w
14922     ##1
14923     \s_keyval_nil
14924     \s_keyval_mark #1 { }
14925     \s_keyval_mark \_keyval_trim_auxii:w
14926     \_keyval_trim_auxiii:w
14927     #1 \s_keyval_nil
14928     \_keyval_trim_auxiv:w
14929   }
14930   \cs_new:Npn \_keyval_trim_auxi:w ##1 \s_keyval_mark #1 ##2 \s_keyval_mark ##3
14931   {
14932     ##3
14933     \_keyval_trim_auxi:w
14934     \s_keyval_mark
14935     ##2
14936     \s_keyval_mark #1 {##1}
14937   }
14938   \cs_new:Npn \_keyval_trim_auxii:w \_keyval_trim_auxi:w \s_keyval_mark \s_keyval_m
14939   {
14940     \_keyval_trim_auxiii:w
14941     ##1
14942   }
14943   \cs_new:Npn \_keyval_trim_auxiii:w ##1 #1 \s_keyval_nil ##2
```

```

14944     {
14945         ##2
14946         ##1 \s__keyval_nil
14947         \__keyval_trim_auxiii:w
14948     }

```

This is the one macro which differs from the original definition.

```

14949     \cs_new:Npn \__keyval_trim_auxiv:w
14950         \s__keyval_mark ##1 \s__keyval_nil
14951         \__keyval_trim_auxiii:w \s__keyval_nil \__keyval_trim_auxiii:w
14952         ##2
14953         { ##2 { ##1 } }
14954     }
14955     \__keyval_tmp:n { ~ }
14956 \group_end:

```

(End definition for `__keyval_trim:nN` and others.)

22.2 Constants and variables

```

14957 <@@=keys>

```

Various storage areas for the different data which make up keys.

```

\c__keys_code_root_str
\c__keys_default_root_str
\c__keys_groups_root_str
\c__keys_inherit_root_str
\c__keys_type_root_str
\c__keys_validate_root_str
14958 \str_const:Nn \c__keys_code_root_str { key-code~>~ }
14959 \str_const:Nn \c__keys_default_root_str { key-default~>~ }
14960 \str_const:Nn \c__keys_groups_root_str { key-groups~>~ }
14961 \str_const:Nn \c__keys_inherit_root_str { key-inherit~>~ }
14962 \str_const:Nn \c__keys_type_root_str { key-type~>~ }
14963 \str_const:Nn \c__keys_validate_root_str { key-validate~>~ }

```

(End definition for `\c__keys_code_root_str` and others.)

`\c__keys_props_root_str` The prefix for storing properties.

```

14964 \str_const:Nn \c__keys_props_root_str { key-prop~>~ }

```

(End definition for `\c__keys_props_root_str`.)

`\l_keys_choice_int` `\l_keys_choice_tl` Publicly accessible data on which choice is being used when several are generated as a set.

```

14965 \int_new:N \l_keys_choice_int
14966 \tl_new:N \l_keys_choice_tl

```

(End definition for `\l_keys_choice_int` and `\l_keys_choice_tl`. These variables are documented on page 193.)

`\l__keys_groups_clist` Used for storing and recovering the list of groups which apply to a key: set as a comma list but at one point we have to use this for a token list recovery.

```

14967 \clist_new:N \l__keys_groups_clist

```

(End definition for `\l__keys_groups_clist`.)

`\l_keys_key_str` `\l_keys_key_tl` The name of a key itself: needed when setting keys. The `tl` version is deprecated but has to be handled manually.

```

14968 \str_new:N \l_keys_key_str
14969 \tl_new:N \l_keys_key_tl

```

(End definition for `\l_keys_key_str` and `\l_keys_key_tl`. These variables are documented on page 195.)

`\l__keys_module_str` The module for an entire set of keys.

14970 `\str_new:N \l__keys_module_str`

(End definition for `\l__keys_module_str`.)

`\l__keys_no_value_bool` A marker is needed internally to show if only a key or a key plus a value was seen: this is recorded here.

14971 `\bool_new:N \l__keys_no_value_bool`

(End definition for `\l__keys_no_value_bool`.)

`\l__keys_only_known_bool` Used to track if only “known” keys are being set.

14972 `\bool_new:N \l__keys_only_known_bool`

(End definition for `\l__keys_only_known_bool`.)

`\l_keys_path_str` The “path” of the current key is stored here: this is available to the programmer and so is public. The older version is deprecated but has to be handled manually.

`\l_keys_path_tl`

14973 `\str_new:N \l_keys_path_str`

14974 `\tl_new:N \l_keys_path_tl`

(End definition for `\l_keys_path_str` and `\l_keys_path_tl`. These variables are documented on page 195.)

`\l__keys_inherit_str`

14975 `\str_new:N \l__keys_inherit_str`

(End definition for `\l__keys_inherit_str`.)

`\l__keys_relative_tl` The relative path for passing keys back to the user. As this can be explicitly no-value, it must be a token list.

14976 `\tl_new:N \l__keys_relative_tl`

14977 `\tl_set:Nn \l__keys_relative_tl { \q__keys_no_value }`

(End definition for `\l__keys_relative_tl`.)

`\l__keys_property_str` The “property” begin set for a key at definition time is stored here.

14978 `\str_new:N \l__keys_property_str`

(End definition for `\l__keys_property_str`.)

`\l__keys_selective_bool` Two flags for using key groups: one to indicate that “selective” setting is active, a second to specify which type (“opt-in” or “opt-out”).

`\l__keys_filtered_bool`

14979 `\bool_new:N \l__keys_selective_bool`

14980 `\bool_new:N \l__keys_filtered_bool`

(End definition for `\l__keys_selective_bool` and `\l__keys_filtered_bool`.)

`\l__keys_selective_seq` The list of key groups being filtered in or out during selective setting.

14981 `\seq_new:N \l__keys_selective_seq`

(End definition for `\l__keys_selective_seq`.)

`\l__keys_unused_clist` Used when setting only some keys to store those left over.

```
14982 \tl_new:N \l__keys_unused_clist
```

(End definition for `\l__keys_unused_clist`.)

`\l__keys_value_tl` The value given for a key: may be empty if no value was given.

```
14983 \tl_new:N \l__keys_value_tl
```

(End definition for `\l__keys_value_tl`. This variable is documented on page 195.)

`\l__keys_tmp_bool` Scratch space.

```
14984 \bool_new:N \l__keys_tmp_bool
```

```
14985 \tl_new:N \l__keys_tmpa_tl
```

```
14986 \tl_new:N \l__keys_tmpb_tl
```

(End definition for `\l__keys_tmp_bool`, `\l__keys_tmpa_tl`, and `\l__keys_tmpb_tl`.)

22.2.1 Internal auxiliaries

`\s__keys_nil` Internal scan marks.

```
14987 \scan_new:N \s__keys_nil
```

`\s__keys_mark`

```
14988 \scan_new:N \s__keys_mark
```

`\s__keys_stop`

```
14989 \scan_new:N \s__keys_stop
```

(End definition for `\s__keys_nil`, `\s__keys_mark`, and `\s__keys_stop`.)

`\q__keys_no_value` Internal quarks.

```
14990 \quark_new:N \q__keys_no_value
```

(End definition for `\q__keys_no_value`.)

`__keys_quark_if_no_value_p:N` Branching quark conditional.

```
14991 \__kernel_quark_new_conditional:Nn \__keys_quark_if_no_value:N { TF }
```

`__keys_quark_if_no_value:NTF`

(End definition for `__keys_quark_if_no_value:NTF`.)

22.3 The key defining mechanism

`\keys_define:nn` The public function for definitions is just a wrapper for the lower level mechanism, more or less. The outer function is designed to keep a track of the current module, to allow safe nesting. The module is set removing any leading / (which is not needed here).

```
14992 \cs_new_protected:Npn \keys_define:nn
```

```
14993 { \__keys_define:onnn \l__keys_module_str }
```

```
14994 \cs_new_protected:Npn \__keys_define:nnn #1#2#3
```

```
14995 {
```

```
14996   \str_set:Nx \l__keys_module_str { \__keys_trim_spaces:n {#2} }
```

```
14997   \keyval_parse:NNn \__keys_define:n \__keys_define:nn {#3}
```

```
14998   \str_set:Nn \l__keys_module_str {#1}
```

```
14999 }
```

```
15000 \cs_generate_variant:Nn \__keys_define:nnn { o }
```

(End definition for `\keys_define:nn` and `__keys_define:nnn`. This function is documented on page 188.)

`__keys_define:n` The outer functions here record whether a value was given and then converge on a
`__keys_define:nn` common internal mechanism. There is first a search for a property in the current key
`__keys_define_aux:nn` name, then a check to make sure it is known before the code hands off to the next step.

```

15001 \cs_new_protected:Npn \__keys_define:n #1
15002 {
15003     \bool_set_true:N \l__keys_no_value_bool
15004     \__keys_define_aux:nn {#1} { }
15005 }
15006 \cs_new_protected:Npn \__keys_define:nn #1#2
15007 {
15008     \bool_set_false:N \l__keys_no_value_bool
15009     \__keys_define_aux:nn {#1} {#2}
15010 }
15011 \cs_new_protected:Npn \__keys_define_aux:nn #1#2
15012 {
15013     \__keys_property_find:n {#1}
15014     \cs_if_exist:cTF { \c__keys_props_root_str \l__keys_property_str }
15015     { \__keys_define_code:n {#2} }
15016     {
15017         \str_if_empty:NF \l__keys_property_str
15018         {
15019             \__kernel_msg_error:nnxx { kernel } { key-property-unknown }
15020             \l__keys_property_str \l_keys_path_str
15021         }
15022     }
15023 }

```

(End definition for `__keys_define:n`, `__keys_define:nn`, and `__keys_define_aux:nn`.)

`__keys_property_find:n` Searching for a property means finding the last . in the input, and storing the text before
`__keys_property_find_auxi:w` and after it. Everything is turned into strings, so there is no problem using an x-type
`__keys_property_find_auxii:w` expansion. Since `__keys_trim_spaces:n` will turn its argument into a string anyway,
`__keys_property_find_auxiii:w` this function uses `\cs_set_nopar:Npx` instead of `\tl_set:Nx` to gain some speed.
`__keys_property_find_auxiv:w`

```

15024 \cs_new_protected:Npn \__keys_property_find:n #1
15025 {
15026     \cs_set_nopar:Npx \l__keys_property_str { \__keys_trim_spaces:n { #1 } }
15027     \exp_after:wN \__keys_property_find_auxi:w \l__keys_property_str
15028     \s__keys_nil \__keys_property_find_auxii:w
15029     . \s__keys_nil \__keys_property_find_err:w
15030 }
15031 \cs_new_protected:Npn \__keys_property_find_auxi:w #1 . #2 \s__keys_nil #3
15032 {
15033     #3 #1 \s__keys_mark #2 \s__keys_nil #3
15034 }
15035 \cs_new_protected:Npn \__keys_property_find_auxii:w
15036     #1 \s__keys_mark #2 \s__keys_nil \__keys_property_find_auxii:w . \s__keys_nil
15037     \__keys_property_find_err:w
15038 {
15039     \cs_set_nopar:Npx \l_keys_path_str
15040     { \str_if_empty:NF \l__keys_module_str { \l__keys_module_str / } #1 }
15041     \__keys_property_find_auxi:w #2 \s__keys_nil \__keys_property_find_auxiii:w . \s__keys_
15042     \__keys_property_find_auxiv:w
15043 }
15044 \cs_new_protected:Npn \__keys_property_find_auxiii:w #1 \s__keys_mark

```



```

15045 {
15046   \cs_set_nopar:Npx \l_keys_path_str { \l_keys_path_str . #1 }
15047   \__keys_property_find_auxi:w
15048 }
15049 \cs_new_protected:Npn \__keys_property_find_auxiv:w
15050   #1 \s_keys_nil \__keys_property_find_auxiii:w
15051   \s_keys_mark \s_keys_nil \__keys_property_find_auxiv:w
15052 {
15053   \cs_set_nopar:Npx \l__keys_property_str { . #1 }
15054   \cs_set_nopar:Npx \l_keys_path_str
15055     { \exp_after:wN \__keys_trim_spaces:n \exp_after:wN { \l_keys_path_str } }
15056   \tl_set_eq:NN \l_keys_path_tl \l_keys_path_str
15057 }
15058 \cs_new_protected:Npn \__keys_property_find_err:w
15059   #1 \s_keys_nil #2 \__keys_property_find_err:w
15060 {
15061   \str_clear:N \l__keys_property_str
15062   \__kernel_msg_error:nnn { kernel } { key-no-property } {#1}
15063 }

```

(End definition for `__keys_property_find:n` and others.)

`__keys_define_code:n` Two possible cases. If there is a value for the key, then just use the function. If not, then
`__keys_define_code:w` a check to make sure there is no need for a value with the property. If there should be
one then complain, otherwise execute it. There is no need to check for a `:` as if it was
missing the earlier tests would have failed.

```

15064 \cs_new_protected:Npn \__keys_define_code:n #1
15065 {
15066   \bool_if:NTF \l_keys_no_value_bool
15067   {
15068     \exp_after:wN \__keys_define_code:w
15069     \l__keys_property_str \s_keys_stop
15070     { \use:c { \c_keys_props_root_str \l_keys_property_str } }
15071     {
15072       \__kernel_msg_error:nnxx { kernel } { key-property-requires-value }
15073       \l__keys_property_str \l_keys_path_str
15074     }
15075   }
15076   { \use:c { \c_keys_props_root_str \l_keys_property_str } {#1} }
15077 }
15078 \exp_last_unbraced:NNNNo
15079 \cs_new:Npn \__keys_define_code:w #1 \c_colon_str #2 \s_keys_stop
15080   { \tl_if_empty:nTF {#2} }

```

(End definition for `__keys_define_code:n` and `__keys_define_code:w`.)

22.4 Turning properties into actions

`__keys_bool_set:Nn` Boolean keys are really just choices, but all done by hand. The second argument here is
`__keys_bool_set:cn` the scope: either empty or `g` for global.

```

15081 \cs_new_protected:Npn \__keys_bool_set:Nn #1#2
15082 {
15083   \bool_if_exist:NF #1 { \bool_new:N #1 }
15084   \__keys_choice_make:

```

```

15085     \__keys_cmd_set:nx { \l_keys_path_str / true }
15086     { \exp_not:c { bool_ #2 set_true:N } \exp_not:N #1 }
15087     \__keys_cmd_set:nx { \l_keys_path_str / false }
15088     { \exp_not:c { bool_ #2 set_false:N } \exp_not:N #1 }
15089     \__keys_cmd_set:nn { \l_keys_path_str / unknown }
15090     {
15091         \__kernel_msg_error:nxx { kernel } { boolean-values-only }
15092         \l_keys_key_str
15093     }
15094     \__keys_default_set:n { true }
15095 }
15096 \cs_generate_variant:Nn \__keys_bool_set:Nn { c }

```

(End definition for __keys_bool_set:Nn.)

__keys_bool_set_inverse:Nn Inverse boolean setting is much the same.

```

\__keys_bool_set_inverse:cn
15097 \cs_new_protected:Npn \__keys_bool_set_inverse:Nn #1#2
15098 {
15099     \bool_if_exist:NF #1 { \bool_new:N #1 }
15100     \__keys_choice_make:
15101     \__keys_cmd_set:nx { \l_keys_path_str / true }
15102     { \exp_not:c { bool_ #2 set_false:N } \exp_not:N #1 }
15103     \__keys_cmd_set:nx { \l_keys_path_str / false }
15104     { \exp_not:c { bool_ #2 set_true:N } \exp_not:N #1 }
15105     \__keys_cmd_set:nn { \l_keys_path_str / unknown }
15106     {
15107         \__kernel_msg_error:nxx { kernel } { boolean-values-only }
15108         \l_keys_key_str
15109     }
15110     \__keys_default_set:n { true }
15111 }
15112 \cs_generate_variant:Nn \__keys_bool_set_inverse:Nn { c }

```

(End definition for __keys_bool_set_inverse:Nn.)

__keys_choice_make: To make a choice from a key, two steps: set the code, and set the unknown key. As
 __keys_multichoice_make: multichoicees and choices are essentially the same bar one function, the code is given
 __keys_choice_make:N together.
 __keys_choice_make_aux:N

```

15113 \cs_new_protected:Npn \__keys_choice_make:
15114 { \__keys_choice_make:N \__keys_choice_find:n }
15115 \cs_new_protected:Npn \__keys_multichoice_make:
15116 { \__keys_choice_make:N \__keys_multichoice_find:n }
15117 \cs_new_protected:Npn \__keys_choice_make:N #1
15118 {
15119     \cs_if_exist:cTF
15120     { \c__keys_type_root_str \__keys_parent:o \l_keys_path_str }
15121     {
15122         \str_if_eq:vnTF
15123         { \c__keys_type_root_str \__keys_parent:o \l_keys_path_str }
15124         { choice }
15125         {
15126             \__kernel_msg_error:nxxx { kernel } { nested-choice-key }
15127             \l_keys_path_tl { \__keys_parent:o \l_keys_path_str }
15128         }

```

```

15129         { \__keys_choice_make_aux:N #1 }
15130     }
15131     { \__keys_choice_make_aux:N #1 }
15132 }
15133 \cs_new_protected:Npn \__keys_choice_make_aux:N #1
15134 {
15135     \cs_set_nopar:cpn { \c__keys_type_root_str \l_keys_path_str }
15136     { choice }
15137     \__keys_cmd_set:nn \l_keys_path_str { #1 {##1} }
15138     \__keys_cmd_set:nn { \l_keys_path_str / unknown }
15139     {
15140         \__kernel_msg_error:nnxx { kernel } { key-choice-unknown }
15141         \l_keys_path_str {##1}
15142     }
15143 }

```

(End definition for __keys_choice_make: and others.)

__keys_choices_make:nn Auto-generating choices means setting up the root key as a choice, then defining each
 __keys_multichoices_make:nn choice in turn.

```

\__keys_choices_make:Nnn
15144 \cs_new_protected:Npn \__keys_choices_make:nn
15145 { \__keys_choices_make:Nnn \__keys_choice_make: }
15146 \cs_new_protected:Npn \__keys_multichoices_make:nn
15147 { \__keys_choices_make:Nnn \__keys_multichoice_make: }
15148 \cs_new_protected:Npn \__keys_choices_make:Nnn #1#2#3
15149 {
15150     #1
15151     \int_zero:N \l_keys_choice_int
15152     \clist_map_inline:nn {#2}
15153     {
15154         \int_incr:N \l_keys_choice_int
15155         \__keys_cmd_set:nx
15156         { \l_keys_path_str / \__keys_trim_spaces:n {##1} }
15157         {
15158             \tl_set:Nn \exp_not:N \l_keys_choice_tl {##1}
15159             \int_set:Nn \exp_not:N \l_keys_choice_int
15160             { \int_use:N \l_keys_choice_int }
15161             \exp_not:n {#3}
15162         }
15163     }
15164 }

```

(End definition for __keys_choices_make:nn, __keys_multichoices_make:nn, and __keys_choices_make:Nnn.)

__keys_cmd_set:nn Setting the code for a key first logs if appropriate that we are defining a new key, then
 __keys_cmd_set:nx saves the code.

```

\__keys_cmd_set:Vn
\__keys_cmd_set:Vo
15165 \cs_new_protected:Npn \__keys_cmd_set:nn #1#2
15166 { \cs_set_protected:cpn { \c__keys_code_root_str #1 } ##1 {#2} }
15167 \cs_generate_variant:Nn \__keys_cmd_set:nn { nx , Vn , Vo }

```

(End definition for __keys_cmd_set:nn.)

`__keys_cs_set:NNpn` Creating control sequences is a bit more tricky than other cases as we need to pick up the `p` argument. To make the internals look clearer, the trailing `n` argument here is just for appearance.

```

15168 \cs_new_protected:Npn \__keys_cs_set:NNpn #1#2#3#
15169 {
15170   \cs_set_protected:cpx { \c__keys_code_root_str \l_keys_path_str } ##1
15171   { #1 \exp_not:N #2 \exp_not:n {#3} {##1} }
15172   \use_none:n
15173 }
15174 \cs_generate_variant:Nn \__keys_cs_set:NNpn { Nc }

```

(End definition for __keys_cs_set:NNpn.)

`__keys_default_set:n` Setting a default value is easy. These are stored using `\cs_set:cpx` as this avoids any worries about whether a token list exists.

```

15175 \cs_new_protected:Npn \__keys_default_set:n #1
15176 {
15177   \tl_if_empty:nTF {#1}
15178   {
15179     \cs_set_eq:cN
15180     { \c__keys_default_root_str \l_keys_path_str }
15181     \tex_undefined:D
15182   }
15183   {
15184     \cs_set_nopar:cpx
15185     { \c__keys_default_root_str \l_keys_path_str }
15186     { \exp_not:n {#1} }
15187     \__keys_value_requirement:nn { required } { false }
15188   }
15189 }

```

(End definition for __keys_default_set:n.)

`__keys_groups_set:n` Assigning a key to one or more groups uses comma lists. As the list of groups only exists if there is anything to do, the setting is done using a scratch list. For the usual grouping reasons we use the low-level approach to undefining a list. We also use the low-level approach for the other case to avoid tripping up the check-declarations code.

```

15190 \cs_new_protected:Npn \__keys_groups_set:n #1
15191 {
15192   \clist_set:Nn \l__keys_groups_clist {#1}
15193   \clist_if_empty:NTF \l__keys_groups_clist
15194   {
15195     \cs_set_eq:cN { \c__keys_groups_root_str \l_keys_path_str }
15196     \tex_undefined:D
15197   }
15198   {
15199     \cs_set_eq:cN { \c__keys_groups_root_str \l_keys_path_str }
15200     \l__keys_groups_clist
15201   }
15202 }

```

(End definition for __keys_groups_set:n.)

`__keys_inherit:n` Inheritance means ignoring anything already said about the key: zap the lot and set up.

```

15203 \cs_new_protected:Npn \__keys_inherit:n #1
15204 {
15205     \__keys_undefine:
15206     \cs_set_nopar:cpn { \c__keys_inherit_root_str \l_keys_path_str } {#1}
15207 }

```

(End definition for __keys_inherit:n.)

`__keys_initialise:n` A set up for initialisation: just run the code if it exists.

```

15208 \cs_new_protected:Npn \__keys_initialise:n #1
15209 {
15210     \cs_if_exist:cTF
15211     { \c__keys_inherit_root_str \__keys_parent:o \l_keys_path_str }
15212     { \__keys_execute_inherit: }
15213     {
15214         \str_clear:N \l__keys_inherit_str
15215         \cs_if_exist:cT { \c__keys_code_root_str \l_keys_path_str }
15216         { \__keys_execute:nn \l_keys_path_str {#1} }
15217     }
15218 }

```

(End definition for __keys_initialise:n.)

`__keys_meta_make:n` To create a meta-key, simply set up to pass data through.

`__keys_meta_make:nn`

```

15219 \cs_new_protected:Npn \__keys_meta_make:n #1
15220 {
15221     \__keys_cmd_set:Vo \l_keys_path_str
15222     {
15223         \exp_after:wN \keys_set:nn \exp_after:wN { \l__keys_module_str } {#1}
15224     }
15225 }
15226 \cs_new_protected:Npn \__keys_meta_make:nn #1#2
15227 { \__keys_cmd_set:Vn \l_keys_path_str { \keys_set:nn {#1} {#2} } }

```

(End definition for __keys_meta_make:n and __keys_meta_make:nn.)

`__keys_prop_put:Nn` Much the same as other variables, but needs a dedicated auxiliary.

`__keys_prop_put:cn`

```

15228 \cs_new_protected:Npn \__keys_prop_put:Nn #1#2
15229 {
15230     \prop_if_exist:NF #1 { \prop_new:N #1 }
15231     \exp_after:wN \__keys_find_key_module:wNN \l_keys_path_str \s__keys_stop
15232     \l__keys_tmpa_tl \l__keys_tmpb_tl
15233     \__keys_cmd_set:nx \l_keys_path_str
15234     {
15235         \exp_not:c { prop_ #2 put:Nnn }
15236         \exp_not:N #1
15237         { \l__keys_tmpb_tl }
15238         \exp_not:n { {##1} }
15239     }
15240 }
15241 \cs_generate_variant:Nn \__keys_prop_put:Nn { c }

```

(End definition for __keys_prop_put:Nn.)

`__keys_undefine:` Undefining a key has to be done without `\cs_undefine:c` as that function acts globally.

```

15242 \cs_new_protected:Npn \__keys_undefine:
15243 {
15244   \clist_map_inline:nn
15245     { code , default , groups , inherit , type , validate }
15246     {
15247       \cs_set_eq:cN
15248       { \tl_use:c { c__keys_ ##1 _root_str } \l_keys_path_str }
15249       \tex_undefined:D
15250     }
15251 }

```

(End definition for __keys_undefine:.)

`__keys_value_requirement:nn` Validating key input is done using a second function which runs before the main key
`__keys_validate_forbidden:` code. Setting that up means setting it equal to a generic stub which does the check. This
`__keys_validate_required:` approach makes the lookup very fast at the cost of one additional csname per key that
needs it. The cleanup here has to know the structure of the following code.

```

15252 \cs_new_protected:Npn \__keys_value_requirement:nn #1#2
15253 {
15254   \str_case:nnF {#2}
15255   {
15256     { true }
15257     {
15258       \cs_set_eq:cc
15259       { \c__keys_validate_root_str \l_keys_path_str }
15260       { __keys_validate_ #1 : }
15261     }
15262     { false }
15263     {
15264       \cs_if_eq:ccT
15265       { \c__keys_validate_root_str \l_keys_path_str }
15266       { __keys_validate_ #1 : }
15267       {
15268         \cs_set_eq:cN
15269         { \c__keys_validate_root_str \l_keys_path_str }
15270         \tex_undefined:D
15271       }
15272     }
15273   }
15274   {
15275     \__kernel_msg_error:nxx { kernel }
15276     { key-property-boolean-values-only }
15277     { .value_ #1 :n }
15278   }
15279 }
15280 \cs_new_protected:Npn \__keys_validate_forbidden:
15281 {
15282   \bool_if:NF \l__keys_no_value_bool
15283   {
15284     \__kernel_msg_error:nxx { kernel } { value-forbidden }
15285     \l_keys_path_str \l_keys_value_tl
15286     \use_none:nnn
15287   }

```

```

15288 }
15289 \cs_new_protected:Npn \__keys_validate_required:
15290 {
15291   \bool_if:NT \l__keys_no_value_bool
15292   {
15293     \__kernel_msg_error:nnx { kernel } { value-required }
15294     \l_keys_path_str
15295     \use_none:nnn
15296   }
15297 }

```

(End definition for __keys_value_requirement:nn, __keys_validate_forbidden:, and __keys_validate_required:.)

```

\__keys_variable_set:NnnN
\__keys_variable_set:cnnN
  \__keys_variable_set_required:NnnN
  \__keys_variable_set_required:cnnN

```

Setting a variable takes the type and scope separately so that it is easy to make a new variable if needed.

```

15298 \cs_new_protected:Npn \__keys_variable_set:NnnN #1#2#3#4
15299 {
15300   \use:c { #2_if_exist:NF } #1 { \use:c { #2_new:N } #1 }
15301   \__keys_cmd_set:nx \l_keys_path_str
15302   {
15303     \exp_not:c { #2 _ #3 set:N #4 }
15304     \exp_not:N #1
15305     \exp_not:n { {#1} }
15306   }
15307 }
15308 \cs_generate_variant:Nn \__keys_variable_set:NnnN { c }
15309 \cs_new_protected:Npn \__keys_variable_set_required:NnnN #1#2#3#4
15310 {
15311   \__keys_variable_set:NnnN #1 {#2} {#3} #4
15312   \__keys_value_requirement:nn { required } { true }
15313 }
15314 \cs_generate_variant:Nn \__keys_variable_set_required:NnnN { c }

```

(End definition for __keys_variable_set:NnnN and __keys_variable_set_required:NnnN.)

22.5 Creating key properties

The key property functions are all wrappers for internal functions, meaning that things stay readable and can also be altered later on.

Importantly, while key properties have “normal” argument specs, the underlying code always supplies one braced argument to these. As such, argument expansion is handled by hand rather than using the standard tools. This shows up particularly for the two-argument properties, where things would otherwise go badly wrong.

```

.bool_set:N
.bool_set:c
.bool_gset:N
.bool_gset:c

```

One function for this.

```

15315 \cs_new_protected:cpn { \c__keys_props_root_str .bool_set:N } #1
15316 { \__keys_bool_set:Nn #1 { } }
15317 \cs_new_protected:cpn { \c__keys_props_root_str .bool_set:c } #1
15318 { \__keys_bool_set:cn {#1} { } }
15319 \cs_new_protected:cpn { \c__keys_props_root_str .bool_gset:N } #1
15320 { \__keys_bool_set:Nn #1 { g } }
15321 \cs_new_protected:cpn { \c__keys_props_root_str .bool_gset:c } #1
15322 { \__keys_bool_set:cn {#1} { g } }

```

(End definition for `.bool_set:N` and `.bool_gset:N`. These functions are documented on page 189.)

`.bool_set_inverse:N` One function for this.

`.bool_set_inverse:c` 15323 \cs_new_protected:cpn { \c__keys_props_root_str .bool_set_inverse:N } #1

`.bool_gset_inverse:N` 15324 { __keys_bool_set_inverse:Nn #1 { } }

`.bool_gset_inverse:c` 15325 \cs_new_protected:cpn { \c__keys_props_root_str .bool_set_inverse:c } #1

15326 { __keys_bool_set_inverse:cn {#1} { } }

15327 \cs_new_protected:cpn { \c__keys_props_root_str .bool_gset_inverse:N } #1

15328 { __keys_bool_set_inverse:Nn #1 { g } }

15329 \cs_new_protected:cpn { \c__keys_props_root_str .bool_gset_inverse:c } #1

15330 { __keys_bool_set_inverse:cn {#1} { g } }

(End definition for `.bool_set_inverse:N` and `.bool_gset_inverse:N`. These functions are documented on page 189.)

`.choice:` Making a choice is handled internally, as it is also needed by `.generate_choices:n`.

15331 \cs_new_protected:cpn { \c__keys_props_root_str .choice: }

15332 { __keys_choice_make: }

(End definition for `.choice:`. This function is documented on page 189.)

`.choices:nn` For auto-generation of a series of mutually-exclusive choices. Here, #1 consists of two separate arguments, hence the slightly odd-looking implementation.

`.choices:Vn` 15333 \cs_new_protected:cpn { \c__keys_props_root_str .choices:nn } #1

`.choices:on` 15334 { __keys_choices_make:nn #1 }

`.choices:xn` 15335 \cs_new_protected:cpn { \c__keys_props_root_str .choices:Vn } #1

15336 { \exp_args:NV __keys_choices_make:nn #1 }

15337 \cs_new_protected:cpn { \c__keys_props_root_str .choices:on } #1

15338 { \exp_args:No __keys_choices_make:nn #1 }

15339 \cs_new_protected:cpn { \c__keys_props_root_str .choices:xn } #1

15340 { \exp_args:Nx __keys_choices_make:nn #1 }

(End definition for `.choices:nn`. This function is documented on page 189.)

`.code:n` Creating code is simply a case of passing through to the underlying `set` function.

15341 \cs_new_protected:cpn { \c__keys_props_root_str .code:n } #1

15342 { __keys_cmd_set:nn \l_keys_path_str {#1} }

(End definition for `.code:n`. This function is documented on page 189.)

`.clist_set:N`

`.clist_set:c` 15343 \cs_new_protected:cpn { \c__keys_props_root_str .clist_set:N } #1

`.clist_gset:N` 15344 { __keys_variable_set:NnnN #1 { clist } { } n }

`.clist_gset:c` 15345 \cs_new_protected:cpn { \c__keys_props_root_str .clist_set:c } #1

15346 { __keys_variable_set:cnnN {#1} { clist } { } n }

15347 \cs_new_protected:cpn { \c__keys_props_root_str .clist_gset:N } #1

15348 { __keys_variable_set:NnnN #1 { clist } { g } n }

15349 \cs_new_protected:cpn { \c__keys_props_root_str .clist_gset:c } #1

15350 { __keys_variable_set:cnnN {#1} { clist } { g } n }

(End definition for `.clist_set:N` and `.clist_gset:N`. These functions are documented on page 189.)


```

.cs_set:Np
.cs_set:cp 15351 \cs_new_protected:cpn { \c__keys_props_root_str .cs_set:Np } #1
.cs_set_protected:Np 15352 { \__keys_cs_set:NNpn \cs_set:Npn #1 { } }
.cs_set_protected:cp 15353 \cs_new_protected:cpn { \c__keys_props_root_str .cs_set:cp } #1
.cs_gset:Np 15354 { \__keys_cs_set:Ncpn \cs_set:Npn #1 { } }
.cs_gset:cp 15355 \cs_new_protected:cpn { \c__keys_props_root_str .cs_set_protected:Np } #1
.cs_gset_protected:Np 15356 { \__keys_cs_set:NNpn \cs_set_protected:Npn #1 { } }
.cs_gset_protected:cp 15357 \cs_new_protected:cpn { \c__keys_props_root_str .cs_set_protected:cp } #1
15358 { \__keys_cs_set:Ncpn \cs_set_protected:Npn #1 { } }
15359 \cs_new_protected:cpn { \c__keys_props_root_str .cs_gset:Np } #1
15360 { \__keys_cs_set:NNpn \cs_gset:Npn #1 { } }
15361 \cs_new_protected:cpn { \c__keys_props_root_str .cs_gset:cp } #1
15362 { \__keys_cs_set:Ncpn \cs_gset:Npn #1 { } }
15363 \cs_new_protected:cpn { \c__keys_props_root_str .cs_gset_protected:Np } #1
15364 { \__keys_cs_set:NNpn \cs_gset_protected:Npn #1 { } }
15365 \cs_new_protected:cpn { \c__keys_props_root_str .cs_gset_protected:cp } #1
15366 { \__keys_cs_set:Ncpn \cs_gset_protected:Npn #1 { } }

```

(End definition for .cs_set:Np and others. These functions are documented on page 189.)

```

.default:n Expansion is left to the internal functions.
.default:V 15367 \cs_new_protected:cpn { \c__keys_props_root_str .default:n } #1
.default:o 15368 { \__keys_default_set:n {#1} }
.default:x 15369 \cs_new_protected:cpn { \c__keys_props_root_str .default:V } #1
15370 { \exp_args:NV \__keys_default_set:n #1 }
15371 \cs_new_protected:cpn { \c__keys_props_root_str .default:o } #1
15372 { \exp_args:No \__keys_default_set:n {#1} }
15373 \cs_new_protected:cpn { \c__keys_props_root_str .default:x } #1
15374 { \exp_args:Nx \__keys_default_set:n {#1} }

```

(End definition for .default:n. This function is documented on page 190.)

```

.dim_set:N Setting a variable is very easy: just pass the data along.
.dim_set:c 15375 \cs_new_protected:cpn { \c__keys_props_root_str .dim_set:N } #1
.dim_gset:N 15376 { \__keys_variable_set_required:NnnN #1 { dim } { } n }
.dim_gset:c 15377 \cs_new_protected:cpn { \c__keys_props_root_str .dim_set:c } #1
15378 { \__keys_variable_set_required:cnnN {#1} { dim } { } n }
15379 \cs_new_protected:cpn { \c__keys_props_root_str .dim_gset:N } #1
15380 { \__keys_variable_set_required:NnnN #1 { dim } { g } n }
15381 \cs_new_protected:cpn { \c__keys_props_root_str .dim_gset:c } #1
15382 { \__keys_variable_set_required:cnnN {#1} { dim } { g } n }

```

(End definition for .dim_set:N and .dim_gset:N. These functions are documented on page 190.)

```

.fp_set:N Setting a variable is very easy: just pass the data along.
.fp_set:c 15383 \cs_new_protected:cpn { \c__keys_props_root_str .fp_set:N } #1
.fp_gset:N 15384 { \__keys_variable_set_required:NnnN #1 { fp } { } n }
.fp_gset:c 15385 \cs_new_protected:cpn { \c__keys_props_root_str .fp_set:c } #1
15386 { \__keys_variable_set_required:cnnN {#1} { fp } { } n }
15387 \cs_new_protected:cpn { \c__keys_props_root_str .fp_gset:N } #1
15388 { \__keys_variable_set_required:NnnN #1 { fp } { g } n }
15389 \cs_new_protected:cpn { \c__keys_props_root_str .fp_gset:c } #1
15390 { \__keys_variable_set_required:cnnN {#1} { fp } { g } n }

```

(End definition for .fp_set:N and .fp_gset:N. These functions are documented on page 190.)

.groups:n A single property to create groups of keys.

```
15391 \cs_new_protected:cpn { \c__keys_props_root_str .groups:n } #1
15392 { \__keys_groups_set:n {#1} }
```

(End definition for .groups:n. This function is documented on page 190.)

.inherit:n Nothing complex: only one variant at the moment!

```
15393 \cs_new_protected:cpn { \c__keys_props_root_str .inherit:n } #1
15394 { \__keys_inherit:n {#1} }
```

(End definition for .inherit:n. This function is documented on page 190.)

.initial:n The standard hand-off approach.

```
.initial:V 15395 \cs_new_protected:cpn { \c__keys_props_root_str .initial:n } #1
.initial:o 15396 { \__keys_initialise:n {#1} }
.initial:x 15397 \cs_new_protected:cpn { \c__keys_props_root_str .initial:V } #1
15398 { \exp_args:NV \__keys_initialise:n #1 }
15399 \cs_new_protected:cpn { \c__keys_props_root_str .initial:o } #1
15400 { \exp_args:No \__keys_initialise:n {#1} }
15401 \cs_new_protected:cpn { \c__keys_props_root_str .initial:x } #1
15402 { \exp_args:Nx \__keys_initialise:n {#1} }
```

(End definition for .initial:n. This function is documented on page 191.)

.int_set:N Setting a variable is very easy: just pass the data along.

```
.int_set:c 15403 \cs_new_protected:cpn { \c__keys_props_root_str .int_set:N } #1
.int_gset:N 15404 { \__keys_variable_set_required:NnnN #1 { int } { } n }
.int_gset:c 15405 \cs_new_protected:cpn { \c__keys_props_root_str .int_set:c } #1
15406 { \__keys_variable_set_required:cnnN {#1} { int } { } n }
15407 \cs_new_protected:cpn { \c__keys_props_root_str .int_gset:N } #1
15408 { \__keys_variable_set_required:NnnN #1 { int } { g } n }
15409 \cs_new_protected:cpn { \c__keys_props_root_str .int_gset:c } #1
15410 { \__keys_variable_set_required:cnnN {#1} { int } { g } n }
```

(End definition for .int_set:N and .int_gset:N. These functions are documented on page 191.)

.meta:n Making a meta is handled internally.

```
15411 \cs_new_protected:cpn { \c__keys_props_root_str .meta:n } #1
15412 { \__keys_meta_make:n {#1} }
```

(End definition for .meta:n. This function is documented on page 191.)

.meta:nn Meta with path: potentially lots of variants, but for the moment no so many defined.

```
15413 \cs_new_protected:cpn { \c__keys_props_root_str .meta:nn } #1
15414 { \__keys_meta_make:nn #1 }
```

(End definition for .meta:nn. This function is documented on page 191.)

.multichoice: The same idea as .choice: and .choices:nn, but where more than one choice is allowed.

```
.multichoices:nn 15415 \cs_new_protected:cpn { \c__keys_props_root_str .multichoice: }
.multichoices:Vn 15416 { \__keys_multichoice_make: }
.multichoices:on 15417 \cs_new_protected:cpn { \c__keys_props_root_str .multichoices:nn } #1
15418 { \__keys_multichoices_make:nn #1 }
.multichoices:xn 15419 \cs_new_protected:cpn { \c__keys_props_root_str .multichoices:Vn } #1
15420 { \exp_args:NV \__keys_multichoices_make:nn #1 }
```

```

15421 \cs_new_protected:cpn { \c__keys_props_root_str .multichoices:on } #1
15422   { \exp_args:No \__keys_multichoices_make:nn #1 }
15423 \cs_new_protected:cpn { \c__keys_props_root_str .multichoices:xn } #1
15424   { \exp_args:Nx \__keys_multichoices_make:nn #1 }

```

(End definition for .multichoice: and .multichoices:nn. These functions are documented on page 191.)

```

.muskip_set:N Setting a variable is very easy: just pass the data along.
.muskip_set:c 15425 \cs_new_protected:cpn { \c__keys_props_root_str .muskip_set:N } #1
               { \__keys_variable_set_required:NnnN #1 { muskip } { } n }
.muskip_gset:N 15426 \cs_new_protected:cpn { \c__keys_props_root_str .muskip_set:c } #1
               { \__keys_variable_set_required:cnnN {#1} { muskip } { } n }
.muskip_gset:c 15427 \cs_new_protected:cpn { \c__keys_props_root_str .muskip_gset:N } #1
               { \__keys_variable_set_required:NnnN #1 { muskip } { g } n }
               15430 \cs_new_protected:cpn { \c__keys_props_root_str .muskip_gset:c } #1
               { \__keys_variable_set_required:cnnN {#1} { muskip } { g } n }

```

(End definition for .muskip_set:N and .muskip_gset:N. These functions are documented on page 191.)

```

.prop_put:N Setting a variable is very easy: just pass the data along.
.prop_put:c 15433 \cs_new_protected:cpn { \c__keys_props_root_str .prop_put:N } #1
               { \__keys_prop_put:Nn #1 { } }
.prop_gput:N 15434 \cs_new_protected:cpn { \c__keys_props_root_str .prop_put:c } #1
               { \__keys_prop_put:cn {#1} { } }
.prop_gput:c 15435 \cs_new_protected:cpn { \c__keys_props_root_str .prop_gput:N } #1
               { \__keys_prop_put:Nn #1 { g } }
               15438 \cs_new_protected:cpn { \c__keys_props_root_str .prop_gput:c } #1
               { \__keys_prop_put:cn {#1} { g } }

```

(End definition for .prop_put:N and .prop_gput:N. These functions are documented on page 191.)

```

.skip_set:N Setting a variable is very easy: just pass the data along.
.skip_set:c 15441 \cs_new_protected:cpn { \c__keys_props_root_str .skip_set:N } #1
               { \__keys_variable_set_required:NnnN #1 { skip } { } n }
.skip_gset:N 15442 \cs_new_protected:cpn { \c__keys_props_root_str .skip_set:c } #1
               { \__keys_variable_set_required:cnnN {#1} { skip } { } n }
.skip_gset:c 15443 \cs_new_protected:cpn { \c__keys_props_root_str .skip_gset:N } #1
               { \__keys_variable_set_required:NnnN #1 { skip } { g } n }
               15446 \cs_new_protected:cpn { \c__keys_props_root_str .skip_gset:c } #1
               { \__keys_variable_set_required:cnnN {#1} { skip } { g } n }

```

(End definition for .skip_set:N and .skip_gset:N. These functions are documented on page 192.)

```

.tl_set:N Setting a variable is very easy: just pass the data along.
.tl_set:c 15449 \cs_new_protected:cpn { \c__keys_props_root_str .tl_set:N } #1
               { \__keys_variable_set:NnnN #1 { tl } { } n }
.tl_gset:N 15450 \cs_new_protected:cpn { \c__keys_props_root_str .tl_set:c } #1
               { \__keys_variable_set:cnnN {#1} { tl } { } n }
.tl_gset:c 15451 \cs_new_protected:cpn { \c__keys_props_root_str .tl_set_x:N } #1
               { \__keys_variable_set:NnnN #1 { tl } { } x }
.tl_set_x:N 15452 \cs_new_protected:cpn { \c__keys_props_root_str .tl_set_x:c } #1
               { \__keys_variable_set:cnnN {#1} { tl } { } x }
.tl_gset_x:N 15453 \cs_new_protected:cpn { \c__keys_props_root_str .tl_gset:N } #1
               { \__keys_variable_set:NnnN #1 { tl } { g } n }
.tl_gset_x:c 15454 \cs_new_protected:cpn { \c__keys_props_root_str .tl_gset_x:N } #1
               { \__keys_variable_set:cnnN {#1} { tl } { g } n }

```

```

15459 \cs_new_protected:cpn { \c__keys_props_root_str .tl_gset:c } #1
15460 { \__keys_variable_set:cnnN {#1} { tl } { g } n }
15461 \cs_new_protected:cpn { \c__keys_props_root_str .tl_gset_x:N } #1
15462 { \__keys_variable_set:NnnN #1 { tl } { g } x }
15463 \cs_new_protected:cpn { \c__keys_props_root_str .tl_gset_x:c } #1
15464 { \__keys_variable_set:cnnN {#1} { tl } { g } x }

```

(End definition for .tl_set:N and others. These functions are documented on page 192.)

.undefine: Another simple wrapper.

```

15465 \cs_new_protected:cpn { \c__keys_props_root_str .undefine: }
15466 { \__keys_undefine: }

```

(End definition for .undefine:. This function is documented on page 192.)

.value_forbidden:n These are very similar, so both call the same function.

```

15467 \cs_new_protected:cpn { \c__keys_props_root_str .value_forbidden:n } #1
15468 { \__keys_value_requirement:nn { forbidden } {#1} }
15469 \cs_new_protected:cpn { \c__keys_props_root_str .value_required:n } #1
15470 { \__keys_value_requirement:nn { required } {#1} }

```

(End definition for .value_forbidden:n and .value_required:n. These functions are documented on page 192.)

22.6 Setting keys

\keys_set:nn A simple wrapper allowing for nesting.

```

\keys_set:nV 15471 \cs_new_protected:Npn \keys_set:nn #1#2
\keys_set:nv 15472 {
\keys_set:no 15473   \use:x
\__keys_set:nn 15474   {
\__keys_set:nnn 15475     \bool_set_false:N \exp_not:N \l__keys_only_known_bool
15476     \bool_set_false:N \exp_not:N \l__keys_filtered_bool
15477     \bool_set_false:N \exp_not:N \l__keys_selective_bool
15478     \tl_set:Nn \exp_not:N \l__keys_relative_tl
15479     { \exp_not:N \q__keys_no_value }
15480     \__keys_set:nn \exp_not:n { {#1} {#2} }
15481     \bool_if:NT \l__keys_only_known_bool
15482     { \bool_set_true:N \exp_not:N \l__keys_only_known_bool }
15483     \bool_if:NT \l__keys_filtered_bool
15484     { \bool_set_true:N \exp_not:N \l__keys_filtered_bool }
15485     \bool_if:NT \l__keys_selective_bool
15486     { \bool_set_true:N \exp_not:N \l__keys_selective_bool }
15487     \tl_set:Nn \exp_not:N \l__keys_relative_tl
15488     { \exp_not:o \l__keys_relative_tl }
15489   }
15490 }
15491 \cs_generate_variant:Nn \keys_set:nn { nV , nv , no }
15492 \cs_new_protected:Npn \__keys_set:nn #1#2
15493 { \exp_args:No \__keys_set:nnn \l__keys_module_str {#1} {#2} }
15494 \cs_new_protected:Npn \__keys_set:nnn #1#2#3
15495 {
15496   \str_set:Nx \l__keys_module_str { \__keys_trim_spaces:n {#2} }
15497   \keyval_parse:Nnn \__keys_set_keyval:n \__keys_set_keyval:nn {#3}
15498   \str_set:Nn \l__keys_module_str {#1}
15499 }

```

(End definition for `\keys_set:nn`, `__keys_set:nn`, and `__keys_set:nnn`. This function is documented on page 195.)

```

\keys_set_known:nnN
\keys_set_known:nVN
\keys_set_known:nvN
\keys_set_known:noN
\keys_set_known:nnnN
\keys_set_known:nVnN
\keys_set_known:nvnN
\keys_set_known:nonN
\__keys_set_known:nnnnN
\keys_set_known:nn
\keys_set_known:nV
\keys_set_known:nv
\keys_set_known:no
\__keys_set_known:nnn

```

Setting known keys simply means setting the appropriate flag, then running the standard code. To allow for nested setting, any existing value of `\l__keys_unused_clist` is saved on the stack and reset afterwards. Note that for speed/simplicity reasons we use a `tl` operation to set the `clist` here!

```

15500 \cs_new_protected:Npn \keys_set_known:nnN #1#2#3
15501 {
15502     \exp_args:No \__keys_set_known:nnnnN
15503     \l__keys_unused_clist \q__keys_no_value {#1} {#2} #3
15504 }
15505 \cs_generate_variant:Nn \keys_set_known:nnN { nV , nv , no }
15506 \cs_new_protected:Npn \keys_set_known:nnnN #1#2#3#4
15507 {
15508     \exp_args:No \__keys_set_known:nnnnN
15509     \l__keys_unused_clist {#3} {#1} {#2} #4
15510 }
15511 \cs_generate_variant:Nn \keys_set_known:nnnN { nV , nv , no }
15512 \cs_new_protected:Npn \__keys_set_known:nnnnN #1#2#3#4#5
15513 {
15514     \clist_clear:N \l__keys_unused_clist
15515     \__keys_set_known:nnn {#2} {#3} {#4}
15516     \__kernel_tl_set:Nx #5 { \exp_not:o \l__keys_unused_clist }
15517     \tl_set:Nn \l__keys_unused_clist {#1}
15518 }
15519 \cs_new_protected:Npn \keys_set_known:nn #1#2
15520 { \__keys_set_known:nnn \q__keys_no_value {#1} {#2} }
15521 \cs_generate_variant:Nn \keys_set_known:nn { nV , nv , no }
15522 \cs_new_protected:Npn \__keys_set_known:nnn #1#2#3
15523 {
15524     \use:x
15525     {
15526         \bool_set_true:N \exp_not:N \l__keys_only_known_bool
15527         \bool_set_false:N \exp_not:N \l__keys_filtered_bool
15528         \bool_set_false:N \exp_not:N \l__keys_selective_bool
15529         \tl_set:Nn \exp_not:N \l__keys_relative_tl { \exp_not:n {#1} }
15530         \__keys_set:nn \exp_not:n { {#2} {#3} }
15531         \bool_if:NF \l__keys_only_known_bool
15532         { \bool_set_false:N \exp_not:N \l__keys_only_known_bool }
15533         \bool_if:NT \l__keys_filtered_bool
15534         { \bool_set_true:N \exp_not:N \l__keys_filtered_bool }
15535         \bool_if:NT \l__keys_selective_bool
15536         { \bool_set_true:N \exp_not:N \l__keys_selective_bool }
15537         \tl_set:Nn \exp_not:N \l__keys_relative_tl
15538         { \exp_not:o \l__keys_relative_tl }
15539     }
15540 }

```

(End definition for `\keys_set_known:nnN` and others. These functions are documented on page 196.)

```

\keys_set_filter:nnnN
\keys_set_filter:nnVN
\keys_set_filter:nnvN
\keys_set_filter:nnoN
\keys_set_filter:nnnnN
\keys_set_filter:nnVnN
\keys_set_filter:nnvnN
\keys_set_filter:nnonN
\__keys_set_filter:nnnnnN
\keys_set_filter:nnn
\keys_set_filter:nnV
\keys_set_filter:nnv
\keys_set_filter:nno

```

The idea of setting keys in a selective manner again uses flags wrapped around the basic code. The comments on `\keys_set_known:nnN` also apply here. We have a bit more shuffling to do to keep everything nestable.

```

15541 \cs_new_protected:Npn \keys_set_filter:nnnN #1#2#3#4
15542 {
15543   \exp_args:No \__keys_set_filter:nnnnN
15544     \l__keys_unused_clist
15545     \q__keys_no_value {#1} {#2} {#3} #4
15546 }
15547 \cs_generate_variant:Nn \keys_set_filter:nnnN { nnV , nnv , nno }
15548 \cs_new_protected:Npn \keys_set_filter:nnnnN #1#2#3#4#5
15549 {
15550   \exp_args:No \__keys_set_filter:nnnnN
15551     \l__keys_unused_clist {#4} {#1} {#2} {#3} #5
15552 }
15553 \cs_generate_variant:Nn \keys_set_filter:nnnnN { nnV , nnv , nno }
15554 \cs_new_protected:Npn \__keys_set_filter:nnnnN #1#2#3#4#5#6
15555 {
15556   \clist_clear:N \l__keys_unused_clist
15557   \__keys_set_filter:nnnn {#2} {#3} {#4} {#5}
15558   \__kernel_tl_set:Nx #6 { \exp_not:o \l__keys_unused_clist }
15559   \tl_set:Nn \l__keys_unused_clist {#1}
15560 }
15561 \cs_new_protected:Npn \keys_set_filter:nnn #1#2#3
15562 { \__keys_set_filter:nnnn \q__keys_no_value {#1} {#2} {#3} }
15563 \cs_generate_variant:Nn \keys_set_filter:nnn { nnV , nnv , nno }
15564 \cs_new_protected:Npn \__keys_set_filter:nnnn #1#2#3#4
15565 {
15566   \use:x
15567   {
15568     \bool_set_false:N \exp_not:N \l__keys_only_known_bool
15569     \bool_set_true:N \exp_not:N \l__keys_filtered_bool
15570     \bool_set_true:N \exp_not:N \l__keys_selective_bool
15571     \tl_set:Nn \exp_not:N \l__keys_relative_tl { \exp_not:n {#1} }
15572     \__keys_set_selective:nnn \exp_not:n { {#2} {#3} {#4} }
15573     \bool_if:NT \l__keys_only_known_bool
15574       { \bool_set_true:N \exp_not:N \l__keys_only_known_bool }
15575     \bool_if:NF \l__keys_filtered_bool
15576       { \bool_set_false:N \exp_not:N \l__keys_filtered_bool }
15577     \bool_if:NF \l__keys_selective_bool
15578       { \bool_set_false:N \exp_not:N \l__keys_selective_bool }
15579     \tl_set:Nn \exp_not:N \l__keys_relative_tl
15580       { \exp_not:o \l__keys_relative_tl }
15581   }
15582 }
15583 \cs_new_protected:Npn \keys_set_groups:nnn #1#2#3
15584 {
15585   \use:x
15586   {
15587     \bool_set_false:N \exp_not:N \l__keys_only_known_bool
15588     \bool_set_false:N \exp_not:N \l__keys_filtered_bool
15589     \bool_set_true:N \exp_not:N \l__keys_selective_bool
15590     \tl_set:Nn \exp_not:N \l__keys_relative_tl
15591       { \exp_not:N \q__keys_no_value }
15592     \__keys_set_selective:nnn \exp_not:n { {#1} {#2} {#3} }
15593     \bool_if:NT \l__keys_only_known_bool
15594       { \bool_set_true:N \exp_not:N \l__keys_only_known_bool }

```

```

15595         \bool_if:NF \l__keys_filtered_bool
15596         { \bool_set_true:N \exp_not:N \l__keys_filtered_bool }
15597     \bool_if:NF \l__keys_selective_bool
15598     { \bool_set_false:N \exp_not:N \l__keys_selective_bool }
15599     \tl_set:Nn \exp_not:N \l__keys_relative_tl
15600     { \exp_not:o \l__keys_relative_tl }
15601 }
15602 }
15603 \cs_generate_variant:Nn \keys_set_groups:nnn { nnV , nnv , nno }
15604 \cs_new_protected:Npn \__keys_set_selective:nnn
15605 { \exp_args:No \__keys_set_selective:nnnn \l__keys_selective_seq }
15606 \cs_new_protected:Npn \__keys_set_selective:nnnn #1#2#3#4
15607 {
15608     \seq_set_from_clist:Nn \l__keys_selective_seq {#3}
15609     \__keys_set:nn {#2} {#4}
15610     \tl_set:Nn \l__keys_selective_seq {#1}
15611 }

```

(End definition for `\keys_set_filter:nnnN` and others. These functions are documented on page 197.)

```

__keys_set_keyval:n
__keys_set_keyval:nn
__keys_set_keyval:nnn
__keys_set_keyval:onn
__keys_find_key_module:wNN
__keys_find_key_module_auxi:Nw
__keys_find_key_module_auxii:Nw
__keys_find_key_module_auxiii:Nn
__keys_find_key_module_auxiv:Nw
__keys_set_selective:

```

A shared system once again. First, set the current path and add a default if needed. There are then checks to see if the a value is required or forbidden. If everything passes, move on to execute the code.

```

15612 \cs_new_protected:Npn \__keys_set_keyval:n #1
15613 {
15614     \bool_set_true:N \l__keys_no_value_bool
15615     \__keys_set_keyval:onn \l__keys_module_str {#1} { }
15616 }
15617 \cs_new_protected:Npn \__keys_set_keyval:nn #1#2
15618 {
15619     \bool_set_false:N \l__keys_no_value_bool
15620     \__keys_set_keyval:onn \l__keys_module_str {#1} {#2}
15621 }

```

The key path here can be fully defined, after which there is a search for the key and module names: the user may have passed them with part of what is actually the module (for our purposes) in the key name. As that happens on a per-key basis, we use the stack approach to restore the module name without a group.

```

15622 \cs_new_protected:Npn \__keys_set_keyval:nnn #1#2#3
15623 {
15624     \__kernel_tl_set:Nx \l__keys_path_str
15625     {
15626         \tl_if_blank:nF {#1}
15627         { #1 / }
15628         \__keys_trim_spaces:n {#2}
15629     }
15630     \str_clear:N \l__keys_module_str
15631     \str_clear:N \l__keys_inherit_str
15632     \exp_after:wN \__keys_find_key_module:wNN \l__keys_path_str \s__keys_stop
15633     \l__keys_module_str \l__keys_key_str
15634     \tl_set_eq:NN \l__keys_key_tl \l__keys_key_str
15635     \__keys_value_or_default:n {#3}
15636     \bool_if:NTF \l__keys_selective_bool
15637     \__keys_set_selective:

```

```

15638     \__keys_execute:
15639     \str_set:Nn \l__keys_module_str {#1}
15640 }
15641 \cs_generate_variant:Nn \__keys_set_keyval:nnn { o }

```

This function uses `\cs_set_nopar:Npx` internally for performance reasons, the argument `#1` is already a string in every usage, so turning it into a string again seems unnecessary.

```

15642 \cs_new_protected:Npn \__keys_find_key_module:wNN #1 \s__keys_stop #2 #3
15643 {
15644     \__keys_find_key_module_auxi:Nw #2 #1 \s__keys_nil \__keys_find_key_module_auxii:Nw
15645     / \s__keys_nil \__keys_find_key_module_auxiv:Nw #3
15646 }
15647 \cs_new_protected:Npn \__keys_find_key_module_auxi:Nw #1 #2 / #3 \s__keys_nil #4
15648 {
15649     #4 #1 #2 \s__keys_mark #3 \s__keys_nil #4
15650 }
15651 \cs_new_protected:Npn \__keys_find_key_module_auxii:Nw
15652     #1 #2 \s__keys_mark #3 \s__keys_nil \__keys_find_key_module_auxii:Nw
15653 {
15654     \cs_set_nopar:Npx #1 { \tl_if_empty:NF #1 { #1 / } #2 }
15655     \__keys_find_key_module_auxi:Nw #1 #3 \s__keys_nil \__keys_find_key_module_auxiii:Nw
15656 }
15657 \cs_new_protected:Npn \__keys_find_key_module_auxiii:Nw #1 #2 \s__keys_mark
15658 {
15659     \cs_set_nopar:Npx #1 { \tl_if_empty:NF #1 { #1 / } #2 }
15660     \__keys_find_key_module_auxi:Nw #1
15661 }
15662 \cs_new_protected:Npn \__keys_find_key_module_auxiv:Nw
15663     #1 #2 \s__keys_nil #3 \s__keys_mark
15664     \s__keys_nil \__keys_find_key_module_auxiv:Nw #4
15665 {
15666     \cs_set_nopar:Npn #4 { #2 }
15667 }

```

If selective setting is active, there are a number of possible sub-cases to consider. The key name may not be known at all or if it is, it may not have any groups assigned. There is then the question of whether the selection is opt-in or opt-out.

```

15668 \cs_new_protected:Npn \__keys_set_selective:
15669 {
15670     \cs_if_exist:cTF { \c__keys_groups_root_str \l_keys_path_str }
15671     {
15672         \clist_set_eq:Nc \l__keys_groups_clist
15673         { \c__keys_groups_root_str \l_keys_path_str }
15674         \__keys_check_groups:
15675     }
15676     {
15677         \bool_if:NTF \l__keys_filtered_bool
15678         \__keys_execute:
15679         \__keys_store_unused:
15680     }
15681 }

```

In the case where selective setting requires a comparison of the list of groups which apply to a key with the list of those which have been set active. That requires two mappings, and again a different outcome depending on whether opt-in or opt-out is set.


```

15682 \cs_new_protected:Npn \__keys_check_groups:
15683 {
15684   \bool_set_false:N \l__keys_tmp_bool
15685   \seq_map_inline:Nn \l__keys_selective_seq
15686   {
15687     \clist_map_inline:Nn \l__keys_groups_clist
15688     {
15689       \str_if_eq:nnT {##1} {####1}
15690       {
15691         \bool_set_true:N \l__keys_tmp_bool
15692         \clist_map_break:n \seq_map_break:
15693       }
15694     }
15695   }
15696   \bool_if:NTF \l__keys_tmp_bool
15697   {
15698     \bool_if:NTF \l__keys_filtered_bool
15699     \__keys_store_unused:
15700     \__keys_execute:
15701   }
15702   {
15703     \bool_if:NTF \l__keys_filtered_bool
15704     \__keys_execute:
15705     \__keys_store_unused:
15706   }
15707 }

```

(End definition for __keys_set_keyval:n and others.)

```

\__keys_value_or_default:n If a value is given, return it as #1, otherwise send a default if available.
\__keys_default_inherit:
15708 \cs_new_protected:Npn \__keys_value_or_default:n #1
15709 {
15710   \bool_if:NTF \l__keys_no_value_bool
15711   {
15712     \cs_if_exist:cTF { \c__keys_default_root_str \l_keys_path_str }
15713     {
15714       \tl_set_eq:Nc
15715       \l_keys_value_tl
15716       { \c__keys_default_root_str \l_keys_path_str }
15717     }
15718     {
15719       \tl_clear:N \l_keys_value_tl
15720       \cs_if_exist:cT
15721       { \c__keys_inherit_root_str \__keys_parent:o \l_keys_path_str }
15722       { \__keys_default_inherit: }
15723     }
15724   }
15725   { \tl_set:Nn \l_keys_value_tl {#1} }
15726 }
15727 \cs_new_protected:Npn \__keys_default_inherit:
15728 {
15729   \clist_map_inline:cn
15730   { \c__keys_inherit_root_str \__keys_parent:o \l_keys_path_str }
15731   {

```

```

15732     \cs_if_exist:cT
15733     { \c__keys_default_root_str ##1 / \l_keys_key_str }
15734     {
15735         \tl_set_eq:Nc
15736         \l_keys_value_tl
15737         { \c__keys_default_root_str ##1 / \l_keys_key_str }
15738         \clist_map_break:
15739     }
15740 }
15741 }

```

(End definition for __keys_value_or_default:n and __keys_default_inherit:.)

__keys_execute: Actually executing a key is done in two parts. First, look for the key itself, then look for the **unknown** key with the same path. If both of these fail, complain. What exactly happens if a key is unknown depends on whether unknown keys are being skipped or if an error should be raised.

```

\__keys_execute:nn
\__keys_execute:no
\__keys_store_unused:
\__keys_store_unused_aux:
15742 \cs_new_protected:Npn \__keys_execute:
15743 {
15744     \cs_if_exist:cTF { \c__keys_code_root_str \l_keys_path_str }
15745     {
15746         \cs_if_exist_use:c { \c__keys_validate_root_str \l_keys_path_str }
15747         \__keys_execute:no \l_keys_path_str \l_keys_value_tl
15748     }
15749     {
15750         \cs_if_exist:cTF
15751         { \c__keys_inherit_root_str \__keys_parent:o \l_keys_path_str }
15752         { \__keys_execute_inherit: }
15753         { \__keys_execute_unknown: }
15754     }
15755 }

```

To deal with the case where there is no hit, we leave __keys_execute_unknown: in the input stream and clean it up using the break function: that avoids needing a boolean.

```

15756 \cs_new_protected:Npn \__keys_execute_inherit:
15757 {
15758     \clist_map_inline:cn
15759     { \c__keys_inherit_root_str \__keys_parent:o \l_keys_path_str }
15760     {
15761         \cs_if_exist:cT
15762         { \c__keys_code_root_str ##1 / \l_keys_key_str }
15763         {
15764             \str_set:Nn \l__keys_inherit_str {##1}
15765             \cs_if_exist_use:c { \c__keys_validate_root_str ##1 / \l_keys_key_str }
15766             \__keys_execute:no { ##1 / \l_keys_key_str } \l_keys_value_tl
15767             \clist_map_break:n \use_none:n
15768         }
15769     }
15770     \__keys_execute_unknown:
15771 }
15772 \cs_new_protected:Npn \__keys_execute_unknown:
15773 {
15774     \bool_if:NTF \l__keys_only_known_bool
15775     { \__keys_store_unused: }

```

```

15776 {
15777   \cs_if_exist:cTF
15778     { \c__keys_code_root_str \l__keys_module_str / unknown }
15779     { \__keys_execute:no { \l__keys_module_str / unknown } \l_keys_value_tl }
15780     {
15781       \__kernel_msg_error:nxxx { kernel } { key-unknown }
15782       \l_keys_path_str \l__keys_module_str
15783     }
15784   }
15785 }

```

A key's code is in the control sequence with csname `\c__keys_code_root_str #1`. We expand it once to get the replacement text (with argument #2) and call `\use:n` with this replacement as its argument. This ensures that any undefined control sequence error in the key's code will lead to an error message of the form `<argument>...<control sequence>` in which one can read the (undefined) `<control sequence>` in full, rather than an error message that starts with the potentially very long key name, which would make the (undefined) `<control sequence>` be truncated or sometimes completely hidden. See <https://github.com/latex3/latex2e/issues/351>.

```

15786 \cs_new:Npn \__keys_execute:nn #1#2
15787 { \__keys_execute:no {#1} { \prg_do_nothing: #2 } }
15788 \cs_new:Npn \__keys_execute:no #1#2
15789 {
15790   \exp_args:NNo \exp_args:No \use:n
15791   {
15792     \cs:w \c__keys_code_root_str #1 \exp_after:wN \cs_end:
15793     \exp_after:wN {#2}
15794   }
15795 }

```

When there is no relative path, things here are easy: just save the key name and value. When we are working with a relative path, first we need to turn it into a string: that can't happen earlier as we need to store `\q__keys_no_value`. Then, use a standard delimited approach to fish out the partial path.

```

15796 \cs_new_protected:Npn \__keys_store_unused:
15797 {
15798   \__keys_quark_if_no_value:NTF \l__keys_relative_tl
15799   {
15800     \clist_put_right:Nx \l__keys_unused_clist
15801     {
15802       \l_keys_key_str
15803       \bool_if:NF \l__keys_no_value_bool
15804       { = { \exp_not:o \l_keys_value_tl } }
15805     }
15806   }
15807   {
15808     \tl_if_empty:NTF \l__keys_relative_tl
15809     {
15810       \clist_put_right:Nx \l__keys_unused_clist
15811       {
15812         \l_keys_path_str
15813         \bool_if:NF \l__keys_no_value_bool
15814         { = { \exp_not:o \l_keys_value_tl } }
15815       }

```

```

15816     }
15817     { \__keys_store_unused_aux: }
15818   }
15819 }
15820 \cs_new_protected:Npn \__keys_store_unused_aux:
15821 {
15822   \__kernel_tl_set:Nx \l__keys_relative_tl
15823   { \exp_args:No \__keys_trim_spaces:n \l__keys_relative_tl }
15824   \use:x
15825   {
15826     \cs_set_protected:Npn \__keys_store_unused:w
15827       ###1 \l__keys_relative_tl /
15828       ###2 \l__keys_relative_tl /
15829       ###3 \s__keys_stop
15830   }
15831   {
15832     \tl_if_blank:nF {##1}
15833     {
15834       \__kernel_msg_error:nxxx { kernel } { bad-relative-key-path }
15835       \l_keys_path_str
15836       \l__keys_relative_tl
15837     }
15838     \clist_put_right:Nx \l__keys_unused_clist
15839     {
15840       \exp_not:n {##2}
15841       \bool_if:NF \l__keys_no_value_bool
15842       { = { \exp_not:o \l_keys_value_tl } }
15843     }
15844   }
15845   \use:x
15846   {
15847     \__keys_store_unused:w \l_keys_path_str
15848     \l__keys_relative_tl / \l__keys_relative_tl /
15849     \s__keys_stop
15850   }
15851 }
15852 \cs_new_protected:Npn \__keys_store_unused:w { }

```

(End definition for __keys_execute: and others.)

__keys_choice_find:n Executing a choice has two parts. First, try the choice given, then if that fails call the
 __keys_choice_find:nn unknown key. That always exists, as it is created when a choice is first made. So there
 __keys_multichoice_find:n is no need for any escape code. For multiple choices, the same code ends up used in a
 mapping.

```

15853 \cs_new:Npn \__keys_choice_find:n #1
15854 {
15855   \str_if_empty:NTF \l__keys_inherit_str
15856   { \__keys_choice_find:nn \l_keys_path_str {#1} }
15857   {
15858     \__keys_choice_find:nn
15859     { \l__keys_inherit_str / \l_keys_key_str } {#1}
15860   }
15861 }
15862 \cs_new:Npn \__keys_choice_find:nn #1#2

```

```

15863 {
15864     \cs_if_exist:cTF { \c__keys_code_root_str #1 / \__keys_trim_spaces:n {#2} }
15865     { \__keys_execute:nn { #1 / \__keys_trim_spaces:n {#2} } {#2} }
15866     { \__keys_execute:nn { #1 / unknown } {#2} }
15867 }
15868 \cs_new:Npn \__keys_multichoice_find:n #1
15869 { \clist_map_function:nN {#1} \__keys_choice_find:n }

```

(End definition for __keys_choice_find:n, __keys_choice_find:nn, and __keys_multichoice_find:n.)

22.7 Utilities

```

\__keys_parent:o Used to strip off the ending part of the key path after the last /.
\__keys_parent_auxi:w 15870 \cs_new:Npn \__keys_parent:o #1
\__keys_parent_auxii:w 15871 {
\__keys_parent_auxiii:n 15872     \exp_after:wN \__keys_parent_auxi:w #1 \q_nil \__keys_parent_auxii:w
\__keys_parent_auxiv:w 15873     / \q_nil \__keys_parent_auxiv:w
15874 }
15875 \cs_new:Npn \__keys_parent_auxi:w #1 / #2 \q_nil #3
15876 {
15877     #3 { #1 } #2 \q_nil #3
15878 }
15879 \cs_new:Npn \__keys_parent_auxii:w #1 #2 \q_nil \__keys_parent_auxii:w
15880 {
15881     #1 \__keys_parent_auxi:w #2 \q_nil \__keys_parent_auxiii:n
15882 }
15883 \cs_new:Npn \__keys_parent_auxiii:n #1
15884 {
15885     / #1 \__keys_parent_auxi:w
15886 }
15887 \cs_new:Npn \__keys_parent_auxiv:w #1 \q_nil \__keys_parent_auxiv:w
15888 {
15889 }

```

(End definition for __keys_parent:o and others.)

__keys_trim_spaces:n Space stripping has to allow for the fact that the key here might have several parts, and spaces need to be stripped from each part. Since the key name is turned into a string groups can't be stripped accidentally and the precautions of \tl_trim_spaces:n aren't necessary, in this case it is much faster to just directly strip spaces around /.

```

15890 \group_begin:
15891     \cs_set:Npn \__keys_tmp:n #1
15892     {
15893         \cs_new:Npn \__keys_trim_spaces:n ##1
15894         {
15895             \exp_after:wN \__keys_trim_spaces_auxi:w \tl_to_str:n { / ##1 } /
15896             \s_keys_nil \__keys_trim_spaces_auxi:w
15897             \s_keys_mark \__keys_trim_spaces_auxii:w
15898             #1 / #1
15899             \s_keys_nil \__keys_trim_spaces_auxii:w
15900             \s_keys_mark \__keys_trim_spaces_auxiii:w
15901         }
15902     }

```

```

15903 \__keys_tmp:n { ~ }
15904 \group_end:
15905 \cs_new:Npn \__keys_trim_spaces_auxi:w #1 ~ / #2 \s__keys_nil #3
15906 {
15907     #3 #1 / #2 \s__keys_nil #3
15908 }
15909 \cs_new:Npn \__keys_trim_spaces_auxii:w #1 / ~ #2 \s__keys_mark #3
15910 {
15911     #3 #1 / #2 \s__keys_mark #3
15912 }
15913 \cs_new:Npn \__keys_trim_spaces_auxiii:w
15914 / #1 /
15915 \s__keys_nil \__keys_trim_spaces_auxi:w
15916 \s__keys_mark \__keys_trim_spaces_auxii:w
15917 /
15918 \s__keys_nil \__keys_trim_spaces_auxii:w
15919 \s__keys_mark \__keys_trim_spaces_auxiii:w
15920 {
15921     #1
15922 }

```

(End definition for __keys_trim_spaces:n and others.)

\keys_if_exist_p:nn A utility for others to see if a key exists.

\keys_if_exist:nnTF

```

15923 \prg_new_conditional:Npnn \keys_if_exist:nn #1#2 { p , T , F , TF }
15924 {
15925     \cs_if_exist:cTF
15926     { \c__keys_code_root_str \__keys_trim_spaces:n { #1 / #2 } }
15927     { \prg_return_true: }
15928     { \prg_return_false: }
15929 }

```

(End definition for \keys_if_exist:nnTF. This function is documented on page 197.)

\keys_if_choice_exist_p:nnn Just an alternative view on \keys_if_exist:nnTF.

\keys_if_choice_exist:nnnTF

```

15930 \prg_new_conditional:Npnn \keys_if_choice_exist:nnn #1#2#3
15931 { p , T , F , TF }
15932 {
15933     \cs_if_exist:cTF
15934     { \c__keys_code_root_str \__keys_trim_spaces:n { #1 / #2 / #3 } }
15935     { \prg_return_true: }
15936     { \prg_return_false: }
15937 }

```

(End definition for \keys_if_choice_exist:nnnTF. This function is documented on page 198.)

\keys_show:nn To show a key, show its code using a message.

\keys_log:nn

```

15938 \cs_new_protected:Npn \keys_show:nn
15939 { \__keys_show:Nnn \msg_show:nnxxxx }
15940 \cs_new_protected:Npn \keys_log:nn
15941 { \__keys_show:Nnn \msg_log:nnxxxx }
15942 \cs_new_protected:Npn \__keys_show:Nnn #1#2#3
15943 {
15944     #1 { LaTeX / kernel } { show-key }
15945     { \__keys_trim_spaces:n { #2 / #3 } }

```

```

15946     {
15947         \keys_if_exist:nnT {#2} {#3}
15948         {
15949             \exp_args:Nnf \msg_show_item_unbraced:nn { code }
15950             {
15951                 \exp_args:Nc \cs_replacement_spec:N
15952                 {
15953                     \c__keys_code_root_str
15954                     \__keys_trim_spaces:n { #2 / #3 }
15955                 }
15956             }
15957         }
15958     }
15959     { } { }
15960 }

```

(End definition for `\keys_show:nn`, `\keys_log:nn`, and `__keys_show:Nnn`. These functions are documented on page 198.)

22.8 Messages

For when there is a need to complain.

```

15961 \__kernel_msg_new:nnnn { kernel } { bad-relative-key-path }
15962 { The-key~'#1'~is-not~inside~the~'#2'~path. }
15963 { The-key~'#1'~cannot-be-expressed-relative-to-path~'#2'. }
15964 \__kernel_msg_new:nnnn { kernel } { boolean-values-only }
15965 { Key~'#1'~accepts~boolean-values-only. }
15966 { The-key~'#1'~only-accepts~the-values~'true'~and~'false'. }
15967 \__kernel_msg_new:nnnn { kernel } { key-choice-unknown }
15968 { Key~'#1'~accepts-only~a~fixed~set~of~choices. }
15969 {
15970     The-key~'#1'~only-accepts~predefined-values,~
15971     and~'#2'~is~not~one~of~these.
15972 }
15973 \__kernel_msg_new:nnnn { kernel } { key-unknown }
15974 { The-key~'#1'~is-unknown~and-is-being-ignored. }
15975 {
15976     The-module~'#2'~does~not~have~a~key~called~'#1'.\\
15977     Check~that~you~have~spelled~the~key~name~correctly.
15978 }
15979 \__kernel_msg_new:nnnn { kernel } { nested-choice-key }
15980 { Attempt-to-define~'#1'~as~a~nested~choice~key. }
15981 {
15982     The-key~'#1'~cannot-be-defined-as~a~choice-as~the~parent~key~'#2'~is~
15983     itself~a~choice.
15984 }
15985 \__kernel_msg_new:nnnn { kernel } { value-forbidden }
15986 { The-key~'#1'~does~not~take~a~value. }
15987 {
15988     The-key~'#1'~should-be-given~without~a~value.\\
15989     The-value~'#2'~was~present:~the-key~will~be~ignored.
15990 }
15991 \__kernel_msg_new:nnnn { kernel } { value-required }
15992 { The-key~'#1'~requires~a~value. }

```

```

15993 {
15994   The~key~'#1'~must~have~a~value.\\
15995   No~value~was~present:~the~key~will~be~ignored.
15996 }
15997 \__kernel_msg_new:nnn { kernel } { show-key }
15998 {
15999   The~key~#1~
16000   \tl_if_empty:nTF {#2}
16001   { is~undefined. }
16002   { has~the~properties: #2 . }
16003 }
16004 \</package>

```

23 l3intarray implementation

```

16005 \*package>
16006 \<@@=intarray>

```

23.1 Allocating arrays

`__intarray_entry:w` We use these primitives quite a lot in this module.

`__intarray_count:w` 16007 \cs_new_eq:NN __intarray_entry:w \tex_fontdimen:D
16008 \cs_new_eq:NN __intarray_count:w \tex_hyphenchar:D

(End definition for __intarray_entry:w and __intarray_count:w.)

`\l__intarray_loop_int` A loop index.

16009 \int_new:N \l__intarray_loop_int

(End definition for \l__intarray_loop_int.)

`\c__intarray_sp_dim` Used to convert integers to dimensions fast.

16010 \dim_const:Nn \c__intarray_sp_dim { 1 sp }

(End definition for \c__intarray_sp_dim.)

`\g__intarray_font_int` Used to assign one font per array.

16011 \int_new:N \g__intarray_font_int

(End definition for \g__intarray_font_int.)

16012 __kernel_msg_new:nnn { kernel } { negative-array-size }
16013 { Size~of~array~may~not~be~negative:~#1 }

\intarray_new:Nn Declare #1 to be a font (arbitrarily cmr10 at a never-used size). Store the array's size as the \hyphenchar of that font and make sure enough \fontdimen are allocated, by setting the last one. Then clear any \fontdimen that cmr10 starts with. It seems LuaTeX's cmr10 has an extra \fontdimen parameter number 8 compared to other engines (for a math font we would replace 8 by 22 or some such). Every intarray must be global; it's enough to run this check in \intarray_new:Nn.

\intarray_new:cn

__intarray_new:N

16014 \cs_new_protected:Npn __intarray_new:N #1
16015 {
16016 __kernel_chk_if_free_cs:N #1
16017 \int_gincr:N \g__intarray_font_int


```

16018 \tex_global:D \tex_font:D #1
16019 = cmr10~at~ \g__intarray_font_int \c__intarray_sp_dim \scan_stop:
16020 \int_step_inline:nn { 8 }
16021 { \__kernel_intarray_gset:Nnn #1 {##1} \c_zero_int }
16022 }
16023 \cs_new_protected:Npn \intarray_new:Nn #1#2
16024 {
16025 \__intarray_new:N #1
16026 \__intarray_count:w #1 = \int_eval:n {#2} \scan_stop:
16027 \int_compare:nNnT { \intarray_count:N #1 } < 0
16028 {
16029 \__kernel_msg_error:nxx { kernel } { negative-array-size }
16030 { \intarray_count:N #1 }
16031 }
16032 \int_compare:nNnT { \intarray_count:N #1 } > 0
16033 { \__kernel_intarray_gset:Nnn #1 { \intarray_count:N #1 } { 0 } }
16034 }
16035 \cs_generate_variant:Nn \intarray_new:Nn { c }

```

(End definition for `\intarray_new:Nn` and `__intarray_new:N`. This function is documented on page 201.)

`\intarray_count:N` Size of an array.

```

\intarray_count:c 16036 \cs_new:Npn \intarray_count:N #1 { \int_value:w \__intarray_count:w #1 }
16037 \cs_generate_variant:Nn \intarray_count:N { c }

```

(End definition for `\intarray_count:N`. This function is documented on page 201.)

23.2 Array items

`__intarray_signed_max_dim:n` Used when an item to be stored is larger than `\c_max_dim` in absolute value; it is replaced by $\pm\c_max_dim$.

```

16038 \cs_new:Npn \__intarray_signed_max_dim:n #1
16039 { \int_value:w \int_compare:nNnT {#1} < 0 { - } \c_max_dim }

```

(End definition for `__intarray_signed_max_dim:n`.)

`__intarray_bounds:NNnTF` The functions `\intarray_gset:Nnn` and `\intarray_item:Nn` share bounds checking. The T branch is used if #3 is within bounds of the array #2.

```

16040 \cs_new:Npn \__intarray_bounds:NNnTF #1#2#3
16041 {
16042 \if_int_compare:w 1 > #3 \exp_stop_f:
16043 \__intarray_bounds_error:NNnw #1 #2 {#3}
16044 \else:
16045 \if_int_compare:w #3 > \intarray_count:N #2 \exp_stop_f:
16046 \__intarray_bounds_error:NNnw #1 #2 {#3}
16047 \fi:
16048 \fi:
16049 \use_i:nn
16050 }
16051 \cs_new:Npn \__intarray_bounds_error:NNnw #1#2#3#4 \use_i:nn #5#6
16052 {
16053 #4
16054 #1 { kernel } { out-of-bounds }

```

```

16055     { \token_to_str:N #2 } {#3} { \intarray_count:N #2 }
16056     #6
16057   }

```

(End definition for `__intarray_bounds:NNnTF` and `__intarray_bounds_error:NNnw`.)

`\intarray_gset:Nnn`

Set the appropriate `\fontdimen`. The `__kernel_intarray_gset:Nnn` function does not use `\int_eval:n`, namely its arguments must be suitable for `\int_value:w`. The user version checks the position and value are within bounds.

`\intarray_gset:cnn`

`__kernel_intarray_gset:Nnn`

`__intarray_gset:Nnn`

`__intarray_gset_overflow:Nnn`

```

16058 \cs_new_protected:Npn \__kernel_intarray_gset:Nnn #1#2#3
16059 { \__intarray_entry:w #2 #1 #3 \c__intarray_sp_dim }
16060 \cs_new_protected:Npn \intarray_gset:Nnn #1#2#3
16061 {
16062   \exp_after:wN \__intarray_gset:Nww
16063   \exp_after:wN #1
16064   \int_value:w \int_eval:n {#2} \exp_after:wN ;
16065   \int_value:w \int_eval:n {#3} ;
16066 }
16067 \cs_generate_variant:Nn \intarray_gset:Nnn { c }
16068 \cs_new_protected:Npn \__intarray_gset:Nww #1#2 ; #3 ;
16069 {
16070   \__intarray_bounds:NNnTF \__kernel_msg_error:nxxxx #1 {#2}
16071   {
16072     \__intarray_gset_overflow_test:nw {#3}
16073     \__kernel_intarray_gset:Nnn #1 {#2} {#3}
16074   }
16075   { }
16076 }
16077 \cs_if_exist:NTF \tex_ifabsnum:D
16078 {
16079   \cs_new_protected:Npn \__intarray_gset_overflow_test:nw #1
16080   {
16081     \tex_ifabsnum:D #1 > \c_max_dim
16082     \exp_after:wN \__intarray_gset_overflow:NNnn
16083     \fi:
16084   }
16085 }
16086 {
16087   \cs_new_protected:Npn \__intarray_gset_overflow_test:nw #1
16088   {
16089     \if_int_compare:w \int_abs:n {#1} > \c_max_dim
16090     \exp_after:wN \__intarray_gset_overflow:NNnn
16091     \fi:
16092   }
16093 }
16094 \cs_new_protected:Npn \__intarray_gset_overflow:NNnn #1#2#3#4
16095 {
16096   \__kernel_msg_error:nxxxx { kernel } { overflow }
16097   { \token_to_str:N #2 } {#3} {#4} { \__intarray_signed_max_dim:n {#4} }
16098   #1 #2 {#3} { \__intarray_signed_max_dim:n {#4} }
16099 }

```

(End definition for `\intarray_gset:Nnn` and others. This function is documented on page 201.)

`\intarray_gzero:N` Set the appropriate `\fontdimen` to zero. No bound checking needed. The `\prg_replicate:nn` possibly uses quite a lot of memory, but this is somewhat comparable to the size of the array, and it is much faster than an `\int_step_inline:nn` loop.

```

16100 \cs_new_protected:Npn \intarray_gzero:N #1
16101 {
16102   \int_zero:N \l__intarray_loop_int
16103   \prg_replicate:nn { \intarray_count:N #1 }
16104   {
16105     \int_incr:N \l__intarray_loop_int
16106     \__intarray_entry:w \l__intarray_loop_int #1 \c_zero_dim
16107   }
16108 }
16109 \cs_generate_variant:Nn \intarray_gzero:N { c }

```

(End definition for `\intarray_gzero:N`. This function is documented on page 201.)

`\intarray_item:Nn` Get the appropriate `\fontdimen` and perform bound checks. The `__kernel_intarray_item:Nn` function omits bound checks and omits `\int_eval:n`, namely its argument must be a T_EX integer suitable for `\int_value:w`.

```

\__kernel_intarray_item:Nn
\__intarray_item:Nn
16110 \cs_new:Npn \__kernel_intarray_item:Nn #1#2
16111 { \int_value:w \__intarray_entry:w #2 #1 }
16112 \cs_new:Npn \intarray_item:Nn #1#2
16113 {
16114   \exp_after:wN \__intarray_item:Nw
16115   \exp_after:wN #1
16116   \int_value:w \int_eval:n {#2} ;
16117 }
16118 \cs_generate_variant:Nn \intarray_item:Nn { c }
16119 \cs_new:Npn \__intarray_item:Nw #1#2 ;
16120 {
16121   \__intarray_bounds:NNnTF \__kernel_msg_expandable_error:nnfff #1 {#2}
16122   { \__kernel_intarray_item:Nn #1 {#2} }
16123   { 0 }
16124 }

```

(End definition for `\intarray_item:Nn`, `__kernel_intarray_item:Nn`, and `__intarray_item:Nn`. This function is documented on page 202.)

`\intarray_rand_item:N` Importantly, `\intarray_item:Nn` only evaluates its argument once.

```

\intarray_rand_item:c
16125 \cs_new:Npn \intarray_rand_item:N #1
16126 { \intarray_item:Nn #1 { \int_rand:n { \intarray_count:N #1 } } }
16127 \cs_generate_variant:Nn \intarray_rand_item:N { c }

```

(End definition for `\intarray_rand_item:N`. This function is documented on page 202.)

23.3 Working with contents of integer arrays

`\intarray_const_from_clist:Nn` Similar to `\intarray_new:Nn` (which we don't use because when debugging is enabled that function checks the variable name starts with `g_`). We make use of the fact that T_EX allows allocation of successive `\fontdimen` as long as no other font has been declared: no need to count the comma list items first. We need the code in `\intarray_gset:Nnn` that checks the item value is not too big, namely `__intarray_gset_overflow_test:nw`, but not the code that checks bounds. At the end, set the size of the intarray.

```

16128 \cs_new_protected:Npn \intarray_const_from_clist:Nn #1#2
16129 {
16130   \__intarray_new:N #1
16131   \int_zero:N \l__intarray_loop_int
16132   \clist_map_inline:nn {#2}
16133   { \exp_args:Nf \__intarray_const_from_clist:nN { \int_eval:n {##1} } } #1 }
16134   \__intarray_count:w #1 \l__intarray_loop_int
16135 }
16136 \cs_generate_variant:Nn \intarray_const_from_clist:Nn { c }
16137 \cs_new_protected:Npn \__intarray_const_from_clist:nN #1#2
16138 {
16139   \int_incr:N \l__intarray_loop_int
16140   \__intarray_gset_overflow_test:nw {#1}
16141   \__kernel_intarray_gset:Nnn #2 \l__intarray_loop_int {#1}
16142 }

```

(End definition for `\intarray_const_from_clist:Nn` and `__intarray_const_from_clist:nN`. This function is documented on page 201.)

`\intarray_to_clist:N` Loop through the array, putting a comma before each item. Remove the leading comma with f-expansion. We also use the auxiliary in `\intarray_show:N` with argument comma, space.
`\intarray_to_clist:c`
`__intarray_to_clist:Nn`
`__intarray_to_clist:w`

```

16143 \cs_new:Npn \intarray_to_clist:N #1 { \__intarray_to_clist:Nn #1 { , } }
16144 \cs_generate_variant:Nn \intarray_to_clist:N { c }
16145 \cs_new:Npn \__intarray_to_clist:Nn #1#2
16146 {
16147   \int_compare:nNf { \intarray_count:N #1 } = \c_zero_int
16148   {
16149     \exp_last_unbraced:Nf \use_none:n
16150     { \__intarray_to_clist:w 1 ; #1 {#2} \prg_break_point: }
16151   }
16152 }
16153 \cs_new:Npn \__intarray_to_clist:w #1 ; #2#3
16154 {
16155   \if_int_compare:w #1 > \__intarray_count:w #2
16156   \prg_break:n
16157   \fi:
16158   #3 \__kernel_intarray_item:Nn #2 {#1}
16159   \exp_after:wN \__intarray_to_clist:w
16160   \int_value:w \int_eval:w #1 + \c_one_int ; #2 {#3}
16161 }

```

(End definition for `\intarray_to_clist:N`, `__intarray_to_clist:Nn`, and `__intarray_to_clist:w`. This function is documented on page 270.)

`__kernel_intarray_range_to_clist:Nnn` Loop through part of the array.

```

\__intarray_range_to_clist:ww
16162 \cs_new:Npn \__kernel_intarray_range_to_clist:Nnn #1#2#3
16163 {
16164   \exp_last_unbraced:Nf \use_none:n
16165   {
16166     \exp_after:wN \__intarray_range_to_clist:ww
16167     \int_value:w \int_eval:w #2 \exp_after:wN ;
16168     \int_value:w \int_eval:w #3 ;
16169     #1 \prg_break_point:
16170   }

```

```

16171 }
16172 \cs_new:Npn \__intarray_range_to_clist:ww #1 ; #2 ; #3
16173 {
16174   \if_int_compare:w #1 > #2 \exp_stop_f:
16175   \prg_break:n
16176   \fi:
16177   , \__kernel_intarray_item:Nn #3 {#1}
16178   \exp_after:wN \__intarray_range_to_clist:ww
16179   \int_value:w \int_eval:w #1 + \c_one_int ; #2 ; #3
16180 }

```

(End definition for __kernel_intarray_range_to_clist:Nnn and __intarray_range_to_clist:ww.)

__kernel_intarray_gset_range_from_clist:Nnn Loop through part of the array.

```

\__intarray_gset_range:Nw
16181 \cs_new_protected:Npn \__kernel_intarray_gset_range_from_clist:Nnn #1#2#3
16182 {
16183   \int_set:Nn \l__intarray_loop_int {#2}
16184   \__intarray_gset_range:Nw #1 #3 , , \prg_break_point:
16185 }
16186 \cs_new_protected:Npn \__intarray_gset_range:Nw #1 #2 ,
16187 {
16188   \if_catcode:w \scan_stop: \tl_to_str:n {#2} \scan_stop:
16189   \prg_break:n
16190   \fi:
16191   \__kernel_intarray_gset:Nnn #1 \l__intarray_loop_int {#2}
16192   \int_incr:N \l__intarray_loop_int
16193   \__intarray_gset_range:Nw #1
16194 }

```

(End definition for __kernel_intarray_gset_range_from_clist:Nnn and __intarray_gset_range:Nw.)

```

\intarray_show:N Convert the list to a comma list (with spaces after each comma)
\intarray_show:c
\intarray_log:N
\intarray_log:c
16195 \cs_new_protected:Npn \intarray_show:N { \__intarray_show:NN \msg_show:nnxxxx }
16196 \cs_generate_variant:Nn \intarray_show:N { c }
16197 \cs_new_protected:Npn \intarray_log:N { \__intarray_show:NN \msg_log:nnxxxx }
16198 \cs_generate_variant:Nn \intarray_log:N { c }
16199 \cs_new_protected:Npn \__intarray_show:NN #1#2
16200 {
16201   \__kernel_chk_defined:NT #2
16202   {
16203     #1 { LaTeX/kernel } { show-intarray }
16204     { \token_to_str:N #2 }
16205     { \intarray_count:N #2 }
16206     { >~ \__intarray_to_clist:Nn #2 { , ~ } }
16207     { }
16208   }
16209 }

```

(End definition for \intarray_show:N and \intarray_log:N. These functions are documented on page 202.)

23.4 Random arrays

We only perform the bounds checks once. This is done by two `__intarray_gset_overflow_test:nw`, with an appropriate empty argument to avoid a spurious “at position #1” part in the error message. Then calculate the number of choices: this is at most $(2^{30} - 1) - (-(2^{30} - 1)) + 1 = 2^{31} - 1$, which just barely does not overflow. For small ranges use `__kernel_randint:n` (making sure to subtract 1 *before* adding the random number to the $\langle min \rangle$, to avoid overflow when $\langle min \rangle$ or $\langle max \rangle$ are $\pm \text{c_max_int}$), otherwise `__kernel_randint:nn`. Finally, if there are no random numbers do not define any of the auxiliaries.

```

16210 \cs_new_protected:Npn \intarray_gset_rand:Nn #1
16211 { \intarray_gset_rand:Nnn #1 { 1 } }
16212 \cs_generate_variant:Nn \intarray_gset_rand:Nn { c }
16213 \sys_if_rand_exist:TF
16214 {
16215   \cs_new_protected:Npn \intarray_gset_rand:Nnn #1#2#3
16216   {
16217     \__intarray_gset_rand:Nff #1
16218     { \int_eval:n {#2} } { \int_eval:n {#3} }
16219   }
16220   \cs_new_protected:Npn \__intarray_gset_rand:Nnn #1#2#3
16221   {
16222     \int_compare:nNnTF {#2} > {#3}
16223     {
16224       \__kernel_msg_expandable_error:nnnn
16225       { kernel } { randint-backward-range } {#2} {#3}
16226       \__intarray_gset_rand:Nnn #1 {#3} {#2}
16227     }
16228     {
16229       \__intarray_gset_overflow_test:nw {#2}
16230       \__intarray_gset_rand_auxi:Nnnn #1 { } {#2} {#3}
16231     }
16232   }
16233   \cs_generate_variant:Nn \__intarray_gset_rand:Nnn { Nff }
16234   \cs_new_protected:Npn \__intarray_gset_rand_auxi:Nnnn #1#2#3#4
16235   {
16236     \__intarray_gset_overflow_test:nw {#4}
16237     \__intarray_gset_rand_auxii:Nnnn #1 { } {#4} {#3}
16238   }
16239   \cs_new_protected:Npn \__intarray_gset_rand_auxii:Nnnn #1#2#3#4
16240   {
16241     \exp_args:NNf \__intarray_gset_rand_auxiii:Nnnn #1
16242     { \int_eval:n { #3 - #4 + 1 } } {#4} {#3}
16243   }
16244   \cs_new_protected:Npn \__intarray_gset_rand_auxiii:Nnnn #1#2#3#4
16245   {
16246     \exp_args:NNf \__intarray_gset_all_same:Nn #1
16247     {
16248       \int_compare:nNnTF {#2} > \c__kernel_randint_max_int
16249       {
16250         \exp_stop_f:
16251         \int_eval:n { \__kernel_randint:nn {#3} {#4} }
16252       }

```

```

16253         {
16254             \exp_stop_f:
16255             \int_eval:n { \__kernel_randint:n {#2} - 1 + #3 }
16256         }
16257     }
16258 }
16259 \cs_new_protected:Npn \__intarray_gset_all_same:Nn #1#2
16260 {
16261     \int_zero:N \l__intarray_loop_int
16262     \prg_replicate:nn { \intarray_count:N #1 }
16263     {
16264         \int_incr:N \l__intarray_loop_int
16265         \__kernel_intarray_gset:Nnn #1 \l__intarray_loop_int {#2}
16266     }
16267 }
16268 }
16269 {
16270     \cs_new_protected:Npn \intarray_gset_rand:Nnn #1#2#3
16271     {
16272         \__kernel_msg_error:nnn { kernel } { fp-no-random }
16273         { \intarray_gset_rand:Nnn #1 {#2} {#3} }
16274     }
16275 }
16276 \cs_generate_variant:Nn \intarray_gset_rand:Nnn { c }

```

(End definition for \intarray_gset_rand:Nn and others. These functions are documented on page 270.)

```

16277 \</package>

```

24 l3fp implementation

Nothing to see here: everything is in the subfiles!

25 l3fp-aux implementation

```

16278 \*package>
16279 \@@=fp>

```

25.1 Access to primitives

__fp_int_eval:w Largely for performance reasons, we need to directly access primitives rather than use \int_eval:n. This happens *a lot*, so we use private names. The same is true for __fp_int_eval_end: \romannumeral, although it is used much less widely.

```

16280 \cs_new_eq:NN \__fp_int_eval:w \tex_numexpr:D
16281 \cs_new_eq:NN \__fp_int_eval_end: \scan_stop:
16282 \cs_new_eq:NN \__fp_int_to_roman:w \tex_romannumeral:D

```

(End definition for __fp_int_eval:w, __fp_int_eval_end:, and __fp_int_to_roman:w.)

25.2 Internal representation

Internally, a floating point number $\langle X \rangle$ is a token list containing

```
\s__fp \__fp_chk:w  $\langle case \rangle$   $\langle sign \rangle$   $\langle body \rangle$  ;
```

Let us explain each piece separately.

Internal floating point numbers are used in expressions, and in this context are subject to **f**-expansion. They must leave a recognizable mark after **f**-expansion, to prevent the floating point number from being re-parsed. Thus, `\s__fp` is simply another name for `\relax`.

When used directly without an accessor function, floating points should produce an error: this is the role of `__fp_chk:w`. We could make floating point variables be protected to prevent them from expanding under **x**-expansion, but it seems more convenient to treat them as a subcase of token list variables.

The (decimal part of the) IEEE-754-2008 standard requires the format to be able to represent special floating point numbers besides the usual positive and negative cases. We distinguish the various possibilities by their $\langle case \rangle$, which is a single digit:

- 0 zeros: `+0` and `-0`,
- 1 “normal” numbers (positive and negative),
- 2 infinities: `+inf` and `-inf`,
- 3 quiet and signalling `nan`.

The $\langle sign \rangle$ is 0 (positive) or 2 (negative), except in the case of `nan`, which have $\langle sign \rangle = 1$. This ensures that changing the $\langle sign \rangle$ digit to $2 - \langle sign \rangle$ is exactly equivalent to changing the sign of the number.

Special floating point numbers have the form

```
\s__fp \__fp_chk:w  $\langle case \rangle$   $\langle sign \rangle$  \s__fp_... ;
```

where `\s__fp_...` is a scan mark carrying information about how the number was formed (useful for debugging).

Normal floating point numbers ($\langle case \rangle = 1$) have the form

```
\s__fp \__fp_chk:w 1  $\langle sign \rangle$  { $\langle exponent \rangle$ } { $\langle X_1 \rangle$ } { $\langle X_2 \rangle$ } { $\langle X_3 \rangle$ } { $\langle X_4 \rangle$ } ;
```

Here, the $\langle exponent \rangle$ is an integer, between -10000 and 10000 . The body consists in four blocks of exactly 4 digits, $0000 \leq \langle X_i \rangle \leq 9999$, and the floating point is

$$(-1)^{\langle sign \rangle / 2} \langle X_1 \rangle \langle X_2 \rangle \langle X_3 \rangle \langle X_4 \rangle \cdot 10^{\langle exponent \rangle - 16}$$

where we have concatenated the 16 digits. Currently, floating point numbers are normalized such that the $\langle exponent \rangle$ is minimal, in other words, $1000 \leq \langle X_1 \rangle \leq 9999$.

Calculations are done in base 10000, *i.e.* one myriad.

Table 3: Internal representation of floating point numbers.

Representation	Meaning
0 0 \s_fp_... ;	Positive zero.
0 2 \s_fp_... ;	Negative zero.
1 0 {\langle exponent\rangle} {\langle X ₁ \rangle} {\langle X ₂ \rangle} {\langle X ₃ \rangle} {\langle X ₄ \rangle} ;	Positive floating point.
1 2 {\langle exponent\rangle} {\langle X ₁ \rangle} {\langle X ₂ \rangle} {\langle X ₃ \rangle} {\langle X ₄ \rangle} ;	Negative floating point.
2 0 \s_fp_... ;	Positive infinity.
2 2 \s_fp_... ;	Negative infinity.
3 1 \s_fp_... ;	Quiet nan.
3 1 \s_fp_... ;	Signalling nan.

25.3 Using arguments and semicolons

_fp_use_none_stop_f:n This function removes an argument (typically a digit) and replaces it by \exp_stop_f:, a marker which stops f-type expansion.

```
16283 \cs_new:Npn \_fp\_use\_none\_stop\_f:n #1 { \exp\_stop\_f: }
```

(End definition for _fp_use_none_stop_f:n.)

_fp_use_s:n Those functions place a semicolon after one or two arguments (typically digits).

```
\_fp\_use\_s:nn
16284 \cs_new:Npn \_fp\_use\_s:n #1 { #1; }
16285 \cs_new:Npn \_fp\_use\_s:nn #1#2 { #1#2; }
```

(End definition for _fp_use_s:n and _fp_use_s:nn.)

_fp_use_none_until_s:w Those functions select specific arguments among a set of arguments delimited by a semicolon.

```
\_fp\_use\_i\_until\_s:nw
\_fp\_use\_ii\_until\_s:nnw
16286 \cs_new:Npn \_fp\_use\_none\_until\_s:w #1; { }
16287 \cs_new:Npn \_fp\_use\_i\_until\_s:nw #1#2; {#1}
16288 \cs_new:Npn \_fp\_use\_ii\_until\_s:nnw #1#2#3; {#2}
```

(End definition for _fp_use_none_until_s:w, _fp_use_i_until_s:nw, and _fp_use_ii_until_s:nnw.)

_fp_reverse_args:Nww Many internal functions take arguments delimited by semicolons, and it is occasionally useful to swap two such arguments.

```
16289 \cs_new:Npn \_fp\_reverse\_args:Nww #1 #2; #3; { #1 #3; #2; }
```

(End definition for _fp_reverse_args:Nww.)

_fp_rrot:www Rotate three arguments delimited by semicolons. This is the inverse (or the square) of the Forth primitive ROT, hence the name.

```
16290 \cs_new:Npn \_fp\_rrot:www #1; #2; #3; { #2; #3; #1; }
```

(End definition for _fp_rrot:www.)

_fp_use_i:ww Many internal functions take arguments delimited by semicolons, and it is occasionally useful to remove one or two such arguments.

```
\_fp\_use\_i:www
16291 \cs_new:Npn \_fp\_use\_i:ww #1; #2; { #1; }
16292 \cs_new:Npn \_fp\_use\_i:www #1; #2; #3; { #1; }
```

(End definition for _fp_use_i:ww and _fp_use_i:www.)

25.4 Constants, and structure of floating points

`__fp_misused:n` This receives a floating point object (floating point number or tuple) and generates an error stating that it was misused. This is called when for instance an `fp` variable is left in the input stream and its contents reach `TEX`'s stomach.

```
16293 \cs_new_protected:Npn \__fp_misused:n #1
16294 { \__kernel_msg_error:nnx { kernel } { misused-fp } { \fp_to_tl:n {#1} } }
```

(End definition for `__fp_misused:n`.)

`\s__fp` Floating points numbers all start with `\s__fp __fp_chk:w`, where `\s__fp` is equal to the `TEX` primitive `\relax`, and `__fp_chk:w` is protected. The rest of the floating point number is made of characters (or `\relax`). This ensures that nothing expands under `f`-expansion, nor under `x`-expansion. However, when typeset, `\s__fp` does nothing, and `__fp_chk:w` is expanded. We define `__fp_chk:w` to produce an error.

```
16295 \scan_new:N \s__fp
16296 \cs_new_protected:Npn \__fp_chk:w #1 ;
16297 { \__fp_misused:n { \s__fp \__fp_chk:w #1 ; } }
```

(End definition for `\s__fp` and `__fp_chk:w`.)

`\s__fp_expr_mark` Aliases of `\tex_relax:D`, used to terminate expressions.

```
\s__fp_expr_stop 16298 \scan_new:N \s__fp_expr_mark
16299 \scan_new:N \s__fp_expr_stop
```

(End definition for `\s__fp_expr_mark` and `\s__fp_expr_stop`.)

`\s__fp_mark` Generic scan marks used throughout the module.

```
\s__fp_stop 16300 \scan_new:N \s__fp_mark
16301 \scan_new:N \s__fp_stop
```

(End definition for `\s__fp_mark` and `\s__fp_stop`.)

`_fp_use_i_delimit_by_s_stop:nw` Functions to gobble up to a scan mark.

```
16302 \cs_new:Npn \_fp_use_i_delimit_by_s_stop:nw #1 #2 \s__fp_stop {#1}
```

(End definition for `_fp_use_i_delimit_by_s_stop:nw`.)

`\s__fp_invalid` A couple of scan marks used to indicate where special floating point numbers come from.

```
\s__fp_underflow 16303 \scan_new:N \s__fp_invalid
\s__fp_overflow 16304 \scan_new:N \s__fp_underflow
\s__fp_division 16305 \scan_new:N \s__fp_overflow
\s__fp_exact 16306 \scan_new:N \s__fp_division
16307 \scan_new:N \s__fp_exact
```

(End definition for `\s__fp_invalid` and others.)

`\c_zero_fp` The special floating points. We define the floating points here as “exact”.

```
\c_minus_zero_fp 16308 \tl_const:Nn \c_zero_fp { \s__fp \__fp_chk:w 0 0 \s__fp_exact ; }
\c_inf_fp 16309 \tl_const:Nn \c_minus_zero_fp { \s__fp \__fp_chk:w 0 2 \s__fp_exact ; }
\c_minus_inf_fp 16310 \tl_const:Nn \c_inf_fp { \s__fp \__fp_chk:w 2 0 \s__fp_exact ; }
\c_nan_fp 16311 \tl_const:Nn \c_minus_inf_fp { \s__fp \__fp_chk:w 2 2 \s__fp_exact ; }
16312 \tl_const:Nn \c_nan_fp { \s__fp \__fp_chk:w 3 1 \s__fp_exact ; }
```

(End definition for `\c_zero_fp` and others. These variables are documented on page 210.)

$\backslash c_fp_prec_int$ $\backslash c_fp_half_prec_int$ $\backslash c_fp_block_int$	<p>The number of digits of floating points.</p> <pre> 16313 \int_const:Nn \c_fp_prec_int { 16 } 16314 \int_const:Nn \c_fp_half_prec_int { 8 } 16315 \int_const:Nn \c_fp_block_int { 4 } </pre> <p>(End definition for $\backslash c_fp_prec_int$, $\backslash c_fp_half_prec_int$, and $\backslash c_fp_block_int$.)</p>
$\backslash c_fp_myriad_int$	<p>Blocks have 4 digits so this integer is useful.</p> <pre> 16316 \int_const:Nn \c_fp_myriad_int { 10000 } </pre> <p>(End definition for $\backslash c_fp_myriad_int$.)</p>
$\backslash c_fp_minus_min_exponent_int$ $\backslash c_fp_max_exponent_int$	<p>Normal floating point numbers have an exponent between $- \text{minus_min_exponent}$ and max_exponent inclusive. Larger numbers are rounded to $\pm\infty$. Smaller numbers are rounded to ± 0. It would be more natural to define a min_exponent with the opposite sign but that would waste one T_EX count.</p> <pre> 16317 \int_const:Nn \c_fp_minus_min_exponent_int { 10000 } 16318 \int_const:Nn \c_fp_max_exponent_int { 10000 } </pre> <p>(End definition for $\backslash c_fp_minus_min_exponent_int$ and $\backslash c_fp_max_exponent_int$.)</p>
$\backslash c_fp_max_exp_exponent_int$	<p>If a number's exponent is larger than that, its exponential overflows/underflows.</p> <pre> 16319 \int_const:Nn \c_fp_max_exp_exponent_int { 5 } </pre> <p>(End definition for $\backslash c_fp_max_exp_exponent_int$.)</p>
$\backslash c_fp_overflowing_fp$	<p>A floating point number that is bigger than all normal floating point numbers. This replaces infinities when converting to formats that do not support infinities.</p> <pre> 16320 \tl_const:Nx \c_fp_overflowing_fp 16321 { 16322 \s_fp _fp_chk:w 1 0 16323 { \int_eval:n { \c_fp_max_exponent_int + 1 } } 16324 {1000} {0000} {0000} {0000} ; 16325 } </pre> <p>(End definition for $\backslash c_fp_overflowing_fp$.)</p>
$\backslash _fp_zero_fp:N$ $\backslash _fp_inf_fp:N$	<p>In case of overflow or underflow, we have to output a zero or infinity with a given sign.</p> <pre> 16326 \cs_new:Npn _fp_zero_fp:N #1 16327 { \s_fp _fp_chk:w 0 #1 \s_fp_underflow ; } 16328 \cs_new:Npn _fp_inf_fp:N #1 16329 { \s_fp _fp_chk:w 2 #1 \s_fp_overflow ; } </pre> <p>(End definition for $\backslash _fp_zero_fp:N$ and $\backslash _fp_inf_fp:N$.)</p>
$\backslash _fp_exponent:w$	<p>For normal numbers, the function expands to the exponent, otherwise to 0. This is used in l3str-format.</p> <pre> 16330 \cs_new:Npn _fp_exponent:w \s_fp _fp_chk:w #1 16331 { 16332 \if_meaning:w 1 #1 16333 \exp_after:wN _fp_use_ii_until_s:nnw 16334 \else: 16335 \exp_after:wN _fp_use_i_until_s:nw 16336 \exp_after:wN 0 16337 \fi: 16338 } </pre>

(End definition for `_fp_exponent:w`.)

`_fp_neg_sign:N` When appearing in an integer expression or after `\int_value:w`, this expands to the sign opposite to #1, namely 0 (positive) is turned to 2 (negative), 1 (nan) to 1, and 2 to 0.

```
16339 \cs_new:Npn \_fp_neg_sign:N #1
16340 { \_fp_int_eval:w 2 - #1 \_fp_int_eval_end: }
```

(End definition for `_fp_neg_sign:N`.)

`_fp_kind:w` Expands to 0 for zeros, 1 for normal floating point numbers, 2 for infinities, 3 for NaN, 4 for tuples.

```
16341 \cs_new:Npn \_fp_kind:w #1
16342 {
16343   \_fp_if_type_fp:NTwFw
16344   #1 \_fp_use_ii_until_s:nnw
16345   \s__fp { \_fp_use_i_until_s:nw 4 }
16346   \s__fp_stop
16347 }
```

(End definition for `_fp_kind:w`.)

25.5 Overflow, underflow, and exact zero

`_fp_sanitize:Nw` `_fp_sanitize:wN` `_fp_sanitize_zero:w` Expects the sign and the exponent in some order, then the significand (which we don't touch). Outputs the corresponding floating point number, possibly underflowed to ± 0 or overflowed to $\pm\infty$. The functions `_fp_underflow:w` and `_fp_overflow:w` are defined in `l3fp-traps`.

```
16348 \cs_new:Npn \_fp_sanitize:Nw #1 #2;
16349 {
16350   \if_case:w
16351     \if_int_compare:w #2 > \c_fp_max_exponent_int 1 ~ \else:
16352     \if_int_compare:w #2 < - \c_fp_minus_min_exponent_int 2 ~ \else:
16353     \if_meaning:w 1 #1 3 ~ \fi: \fi: \fi: 0 ~
16354   \or: \exp_after:wN \_fp_overflow:w
16355   \or: \exp_after:wN \_fp_underflow:w
16356   \or: \exp_after:wN \_fp_sanitize_zero:w
16357   \fi:
16358   \s__fp \_fp_chk:w 1 #1 {#2}
16359 }
16360 \cs_new:Npn \_fp_sanitize:wN #1; #2 { \_fp_sanitize:Nw #2 #1; }
16361 \cs_new:Npn \_fp_sanitize_zero:w \s__fp \_fp_chk:w #1 #2 #3;
16362 { \c_zero_fp }
```

(End definition for `_fp_sanitize:Nw`, `_fp_sanitize:wN`, and `_fp_sanitize_zero:w`.)

25.6 Expanding after a floating point number

`_fp_exp_after_o:w` `_fp_exp_after_f:nw` `_fp_exp_after_f:nw` `_fp_exp_after_o:w` `_fp_exp_after_f:nw` `_fp_exp_after_f:nw` Places *tokens* (empty in the case of `_fp_exp_after_o:w`) between the *floating point* and the following tokens, then hits those tokens with o or f-expansion, and leaves the floating point number unchanged.

We first distinguish normal floating points, which have a significand, from the much simpler special floating points.

```

16363 \cs_new:Npn \__fp_exp_after_o:w \s__fp \__fp_chk:w #1
16364 {
16365   \if_meaning:w 1 #1
16366     \exp_after:wN \__fp_exp_after_normal:nNNw
16367   \else:
16368     \exp_after:wN \__fp_exp_after_special:nNNw
16369   \fi:
16370   { }
16371   #1
16372 }
16373 \cs_new:Npn \__fp_exp_after_f:nw #1 \s__fp \__fp_chk:w #2
16374 {
16375   \if_meaning:w 1 #2
16376     \exp_after:wN \__fp_exp_after_normal:nNNw
16377   \else:
16378     \exp_after:wN \__fp_exp_after_special:nNNw
16379   \fi:
16380   { \exp:w \exp_end_continue_f:w #1 }
16381   #2
16382 }

```

(End definition for __fp_exp_after_o:w and __fp_exp_after_f:nw.)

__fp_exp_after_special:nNNw

__fp_exp_after_special:nNNw {<after>} <case> <sign> <scan mark> ;
Special floating point numbers are easy to jump over since they contain few tokens.

```

16383 \cs_new:Npn \__fp_exp_after_special:nNNw #1#2#3#4;
16384 {
16385   \exp_after:wN \s__fp
16386   \exp_after:wN \__fp_chk:w
16387   \exp_after:wN #2
16388   \exp_after:wN #3
16389   \exp_after:wN #4
16390   \exp_after:wN ;
16391   #1
16392 }

```

(End definition for __fp_exp_after_special:nNNw.)

__fp_exp_after_normal:nNNw

For normal floating point numbers, life is slightly harder, since we have many tokens to jump over. Here it would be slightly better if the digits were not braced but instead were delimited arguments (for instance delimited by ,). That may be changed some day.

```

16393 \cs_new:Npn \__fp_exp_after_normal:nNNw #1 1 #2 #3 #4#5#6#7;
16394 {
16395   \exp_after:wN \__fp_exp_after_normal:Nwwwww
16396   \exp_after:wN #2
16397   \int_value:w #3 \exp_after:wN ;
16398   \int_value:w 1 #4 \exp_after:wN ;
16399   \int_value:w 1 #5 \exp_after:wN ;
16400   \int_value:w 1 #6 \exp_after:wN ;
16401   \int_value:w 1 #7 \exp_after:wN ; #1
16402 }
16403 \cs_new:Npn \__fp_exp_after_normal:Nwwwww

```

```

16404      #1 #2; 1 #3 ; 1 #4 ; 1 #5 ; 1 #6 ;
16405      { \s__fp \__fp_chk:w 1 #1 {#2} {#3} {#4} {#5} {#6} ; }

```

(End definition for __fp_exp_after_normal:nNNw.)

25.7 Other floating point types

\s__fp_tuple Floating point tuples take the form \s__fp_tuple __fp_tuple_chk:w { *<fp 1>* *<fp 2>* ... } ; where each *<fp>* is a floating point number or tuple, hence ends with ; itself. When a tuple is typeset, __fp_tuple_chk:w produces an error, just like usual floating point numbers. Tuples may have zero or one element.

```

16406 \scan_new:N \s__fp_tuple
16407 \cs_new_protected:Npn \__fp_tuple_chk:w #1 ;
16408 { \__fp_misused:n { \s__fp_tuple \__fp_tuple_chk:w #1 ; } }
16409 \tl_const:Nn \c__fp_empty_tuple_fp
16410 { \s__fp_tuple \__fp_tuple_chk:w { } ; }

```

(End definition for \s__fp_tuple, __fp_tuple_chk:w, and \c__fp_empty_tuple_fp.)

__fp_tuple_count:w Count the number of items in a tuple of floating points by counting semicolons. The technique is very similar to \tl_count:n, but with the loop built-in. Checking for the end of the loop is done with the \use_none:n #1 construction.

```

16411 \cs_new:Npn \__fp_array_count:n #1
16412 { \__fp_tuple_count:w \s__fp_tuple \__fp_tuple_chk:w {#1} ; }
16413 \cs_new:Npn \__fp_tuple_count:w \s__fp_tuple \__fp_tuple_chk:w #1 ;
16414 {
16415   \int_value:w \__fp_int_eval:w 0
16416   \__fp_tuple_count_loop:Nw #1 { ? \prg_break: } ;
16417   \prg_break_point:
16418   \__fp_int_eval_end:
16419 }
16420 \cs_new:Npn \__fp_tuple_count_loop:Nw #1#2;
16421 { \use_none:n #1 + 1 \__fp_tuple_count_loop:Nw }

```

(End definition for __fp_tuple_count:w, __fp_array_count:n, and __fp_tuple_count_loop:Nw.)

__fp_if_type_fp:NTwFw Used as __fp_if_type_fp:NTwFw *<marker>* {*<true code>*} \s__fp {*<false code>*} \s__fp_stop, this test whether the *<marker>* is \s__fp or not and runs the appropriate *<code>*. The very unusual syntax is for optimization purposes as that function is used for all floating point operations.

```

16422 \cs_new:Npn \__fp_if_type_fp:NTwFw #1 \s__fp #2 #3 \s__fp_stop {#2}

```

(End definition for __fp_if_type_fp:NTwFw.)

__fp_array_if_all_fp:nTF True if all items are floating point numbers. Used for min.

```

16423 \cs_new:Npn \__fp_array_if_all_fp:nTF #1
16424 {
16425   \__fp_array_if_all_fp_loop:w #1 { \s__fp \prg_break: } ;
16426   \prg_break_point: \use_i:nn
16427 }
16428 \cs_new:Npn \__fp_array_if_all_fp_loop:w #1#2 ;
16429 {
16430   \__fp_if_type_fp:NTwFw
16431   #1 \__fp_array_if_all_fp_loop:w

```

```

16432     \s__fp { \prg_break:n \use_iii:nnn }
16433     \s__fp_stop
16434 }

```

(End definition for _fp_array_if_all_fp:nTF and _fp_array_if_all_fp_loop:w.)

_fp_type_from_scan:N Used as _fp_type_from_scan:N *<token>*. Grabs the pieces of the stringified *<token>* which lies after the first s__fp. If the *<token>* does not contain that string, the result is _?.

```

16435 \cs_new:Npn \_fp_type_from_scan:N #1
16436 {
16437     \_fp_if_type_fp:NTwFw
16438     #1 { }
16439     \s__fp { \_fp_type_from_scan_other:N #1 }
16440     \s__fp_stop
16441 }
16442 \cs_new:Npx \_fp_type_from_scan_other:N #1
16443 {
16444     \exp_not:N \exp_after:wN \exp_not:N \_fp_type_from_scan:w
16445     \exp_not:N \token_to_str:N #1 \s__fp_mark
16446     \tl_to_str:n { s__fp _? } \s__fp_mark \s__fp_stop
16447 }
16448 \exp_last_unbraced:NNNNo
16449 \cs_new:Npn \_fp_type_from_scan:w #1
16450 { \tl_to_str:n { s__fp } } #2 \s__fp_mark #3 \s__fp_stop {#2}

```

(End definition for _fp_type_from_scan:N, _fp_type_from_scan_other:N, and _fp_type_from_scan:w.)

_fp_change_func_type:NNN Arguments are *<type marker>* *<function>* *<recovery>*. This gives the function obtained by placing the type after @@. If the function is not defined then *<recovery>* *<function>* is used instead; however that test is not run when the *<type marker>* is s__fp.

```

16451 \cs_new:Npn \_fp_change_func_type:NNN #1#2#3
16452 {
16453     \_fp_if_type_fp:NTwFw
16454     #1 #2
16455     \s__fp
16456     {
16457         \exp_after:wN \_fp_change_func_type_chk:NNN
16458         \cs:w
16459         __fp \_fp_type_from_scan_other:N #1
16460         \exp_after:wN \_fp_change_func_type_aux:w \token_to_str:N #2
16461         \cs_end:
16462         #2 #3
16463     }
16464     \s__fp_stop
16465 }
16466 \exp_last_unbraced:NNNNo
16467 \cs_new:Npn \_fp_change_func_type_aux:w #1 { \tl_to_str:n { __fp } } { }
16468 \cs_new:Npn \_fp_change_func_type_chk:NNN #1#2#3
16469 {
16470     \if_meaning:w \scan_stop: #1
16471     \exp_after:wN #3 \exp_after:wN #2
16472     \else:

```

```

16473     \exp_after:wN #1
16474     \fi:
16475   }

```

(End definition for _fp_change_func_type:NNN, _fp_change_func_type_aux:w, and _fp_change_func_type_chk:NNN.)

```

\_fp_exp_after_any_f:Nnw
\_fp_exp_after_any_f:nw
  \_fp_exp_after_expr_stop_f:nw

```

The Nnw function simply dispatches to the appropriate _fp_exp_after..._f:nw with “...” (either empty or $\langle type \rangle$) extracted from #1, which should start with \s_fp. If it doesn’t start with \s_fp the function _fp_exp_after_?_f:nw defined in l3fp-parse gives an error; another special $\langle type \rangle$ is stop, useful for loops, see below. The nw function has an important optimization for floating points numbers; it also fetches its type marker #2 from the floating point.

```

16476 \cs_new:Npn \_fp_exp_after_any_f:Nnw #1
16477   { \cs:w \_fp_exp_after \_fp_type_from_scan_other:N #1 _f:nw \cs_end: }
16478 \cs_new:Npn \_fp_exp_after_any_f:nw #1#2
16479   {
16480     \_fp_if_type_fp:NTwFw
16481     #2 \_fp_exp_after_f:nw
16482     \s\_fp { \_fp_exp_after_any_f:Nnw #2 }
16483     \s\_fp_stop
16484     {#1} #2
16485   }
16486 \cs_new_eq:NN \_fp_exp_after_expr_stop_f:nw \use_none:nn

```

(End definition for _fp_exp_after_any_f:Nnw, _fp_exp_after_any_f:nw, and _fp_exp_after_expr_stop_f:nw.)

```

\_fp_exp_after_tuple_o:w
\_fp_exp_after_tuple_f:nw
\_fp_exp_after_array_f:w

```

The loop works by using the n argument of _fp_exp_after_any_f:nw to place the loop macro after the next item in the tuple and expand it.

```

  \_fp_exp_after_array_f:w
  \fp1 ;
  ...
  \fpn ;
  \s\_fp_expr_stop

```

```

16487 \cs_new:Npn \_fp_exp_after_tuple_o:w
16488   { \_fp_exp_after_tuple_f:nw { \exp_after:wN \exp_stop_f: } }
16489 \cs_new:Npn \_fp_exp_after_tuple_f:nw
16490   #1 \s\_fp_tuple \_fp_tuple_chk:w #2 ;
16491   {
16492     \exp_after:wN \s\_fp_tuple
16493     \exp_after:wN \_fp_tuple_chk:w
16494     \exp_after:wN {
16495       \exp:w \exp_end_continue_f:w
16496       \_fp_exp_after_array_f:w #2 \s\_fp_expr_stop
16497     }
16498     \exp_after:wN ;
16499     \exp:w \exp_end_continue_f:w #1
16500   }
16501 \cs_new:Npn \_fp_exp_after_array_f:w
16502   { \_fp_exp_after_any_f:nw { \_fp_exp_after_array_f:w } }

```

(End definition for _fp_exp_after_tuple_o:w, _fp_exp_after_tuple_f:nw, and _fp_exp_after_array_f:w.)

25.8 Packing digits

When a positive integer `#1` is known to be less than 10^8 , the following trick splits it into two blocks of 4 digits, padding with zeros on the left.

```
\cs_new:Npn \pack:NNNNNw #1 #2#3#4#5 #6; { {#2#3#4#5} {#6} }
\exp_after:wN \pack:NNNNNw
  \__fp_int_value:w \__fp_int_eval:w 1 0000 0000 + #1 ;
```

The idea is that adding 10^8 to the number ensures that it has exactly 9 digits, and can then easily find which digits correspond to what position in the number. Of course, this can be modified for any number of digits less or equal to 9 (we are limited by `TeX`'s integers). This method is very heavily relied upon in `l3fp-basics`.

More specifically, the auxiliary inserts `+ #1#2#3#4#5 ; {#6}`, which allows us to compute several blocks of 4 digits in a nested manner, performing carries on the fly. Say we want to compute 12345×66778899 . With simplified names, we would do

```
\exp_after:wN \post_processing:w
\__fp_int_value:w \__fp_int_eval:w - 5 0000
  \exp_after:wN \pack:NNNNNw
    \__fp_int_value:w \__fp_int_eval:w 4 9995 0000
      + 12345 * 6677
    \exp_after:wN \pack:NNNNNw
      \__fp_int_value:w \__fp_int_eval:w 5 0000 0000
        + 12345 * 8899 ;
```

The `\exp_after:wN` triggers `\int_value:w __fp_int_eval:w`, which starts a first computation, whose initial value is $-5\,0000$ (the “leading shift”). In that computation appears an `\exp_after:wN`, which triggers the nested computation `\int_value:w __fp_int_eval:w` with starting value $4\,9995\,0000$ (the “middle shift”). That, in turn, expands `\exp_after:wN` which triggers the third computation. The third computation's value is $5\,0000\,0000 + 12345 \times 8899$, which has 9 digits. Adding $5 \cdot 10^8$ to the product allowed us to know how many digits to expect as long as the numbers to multiply are not too big; it also works to some extent with negative results. The `pack` function puts the last 4 of those 9 digits into a brace group, moves the semi-colon delimiter, and inserts a `+`, which combines the carry with the previous computation. The shifts nicely combine into $5\,0000\,0000/10^4 + 4\,9995\,0000 = 5\,0000\,0000$. As long as the operands are in some range, the result of this second computation has 9 digits. The corresponding `pack` function, expanded after the result is computed, braces the last 4 digits, and leaves `+ <5 digits>` for the initial computation. The “leading shift” cancels the combination of the other shifts, and the `\post_processing:w` takes care of packing the last few digits.

Admittedly, this is quite intricate. It is probably the key in making `l3fp` as fast as other pure `TeX` floating point units despite its increased precision. In fact, this is used so much that we provide different sets of packing functions and shifts, depending on ranges of input.

This set of shifts allows for computations involving results in the range $[-4 \cdot 10^8, 5 \cdot 10^8 - 1]$. Shifted values all have exactly 9 digits.

```
\__fp_pack:NNNNNw
\c__fp_trailing_shift_int
\c__fp_middle_shift_int
\c__fp_leading_shift_int
16503 \int_const:Nn \c__fp_leading_shift_int { - 5 0000 }
16504 \int_const:Nn \c__fp_middle_shift_int { 5 0000 * 9999 }
16505 \int_const:Nn \c__fp_trailing_shift_int { 5 0000 * 10000 }
16506 \cs_new:Npn \__fp_pack:NNNNNw #1 #2#3#4#5 #6; { + #1#2#3#4#5 ; {#6} }
```

(End definition for `_fp_pack:NNNNNw` and others.)

`_fp_pack_big:NNNNNw`
`\c_fp_big_trailing_shift_int`
`\c_fp_big_middle_shift_int`
`\c_fp_big_leading_shift_int`

This set of shifts allows for computations involving results in the range $[-5 \cdot 10^8, 6 \cdot 10^8 - 1]$ (actually a bit more). Shifted values all have exactly 10 digits. Note that the upper bound is due to $\text{T}_{\text{E}}\text{X}$'s limit of $2^{31} - 1$ on integers. The shifts are chosen to be roughly the mid-point of 10^9 and 2^{31} , the two bounds on 10-digit integers in $\text{T}_{\text{E}}\text{X}$.

```
16507 \int_const:Nn \c_fp_big_leading_shift_int { - 15 2374 }
16508 \int_const:Nn \c_fp_big_middle_shift_int { 15 2374 * 9999 }
16509 \int_const:Nn \c_fp_big_trailing_shift_int { 15 2374 * 10000 }
16510 \cs_new:Npn \_fp_pack_big:NNNNNw #1#2 #3#4#5#6 #7;
16511 { + #1#2#3#4#5#6 ; {#7} }
```

(End definition for `_fp_pack_big:NNNNNw` and others.)

`_fp_pack_Bigg:NNNNNw`
`\c_fp_Bigg_trailing_shift_int`
`\c_fp_Bigg_middle_shift_int`
`\c_fp_Bigg_leading_shift_int`

This set of shifts allows for computations with results in the range $[-1 \cdot 10^9, 147483647]$; the end-point is $2^{31} - 1 - 2 \cdot 10^9 \simeq 1.47 \cdot 10^8$. Shifted values all have exactly 10 digits.

```
16512 \int_const:Nn \c_fp_Bigg_leading_shift_int { - 20 0000 }
16513 \int_const:Nn \c_fp_Bigg_middle_shift_int { 20 0000 * 9999 }
16514 \int_const:Nn \c_fp_Bigg_trailing_shift_int { 20 0000 * 10000 }
16515 \cs_new:Npn \_fp_pack_Bigg:NNNNNw #1#2 #3#4#5#6 #7;
16516 { + #1#2#3#4#5#6 ; {#7} }
```

(End definition for `_fp_pack_Bigg:NNNNNw` and others.)

`_fp_pack_twice_four:wNNNNNNNN`

`_fp_pack_twice_four:wNNNNNNNN` *(tokens)* ; $\langle \geq 8 \text{ digits} \rangle$

Grabs two sets of 4 digits and places them before the semi-colon delimiter. Putting several copies of this function before a semicolon packs more digits since each takes the digits packed by the others in its first argument.

```
16517 \cs_new:Npn \_fp_pack_twice_four:wNNNNNNNN #1; #2#3#4#5 #6#7#8#9
16518 { #1 {#2#3#4#5} {#6#7#8#9} ; }
```

(End definition for `_fp_pack_twice_four:wNNNNNNNN`.)

`_fp_pack_eight:wNNNNNNNN`

`_fp_pack_eight:wNNNNNNNN` *(tokens)* ; $\langle \geq 8 \text{ digits} \rangle$

Grabs one set of 8 digits and places them before the semi-colon delimiter as a single group. Putting several copies of this function before a semicolon packs more digits since each takes the digits packed by the others in its first argument.

```
16519 \cs_new:Npn \_fp_pack_eight:wNNNNNNNN #1; #2#3#4#5 #6#7#8#9
16520 { #1 {#2#3#4#5#6#7#8#9} ; }
```

(End definition for `_fp_pack_eight:wNNNNNNNN`.)

`_fp_basics_pack_low:NNNNNw`
`_fp_basics_pack_high:NNNNNw`
`_fp_basics_pack_high_carry:w`

Addition and multiplication of significands are done in two steps: first compute a (more or less) exact result, then round and pack digits in the final (braced) form. These functions take care of the packing, with special attention given to the case where rounding has caused a carry. Since rounding can only shift the final digit by 1, a carry always produces an exact power of 10. Thus, `_fp_basics_pack_high_carry:w` is always followed by four times `{0000}`.

This is used in `l3fp-basics` and `l3fp-extended`.

```
16521 \cs_new:Npn \_fp_basics_pack_low:NNNNNw #1 #2#3#4#5 #6;
16522 { + #1 - 1 ; {#2#3#4#5} {#6} ; }
16523 \cs_new:Npn \_fp_basics_pack_high:NNNNNw #1 #2#3#4#5 #6;
16524 {
```

```

16525     \if_meaning:w 2 #1
16526         \__fp_basics_pack_high_carry:w
16527     \fi:
16528     ; {#2#3#4#5} {#6}
16529 }
16530 \cs_new:Npn \__fp_basics_pack_high_carry:w \fi: ; #1
16531 { \fi: + 1 ; {1000} }

```

(End definition for `__fp_basics_pack_low:NNNNw`, `__fp_basics_pack_high:NNNNw`, and `__fp_basics_pack_high_carry:w`.)

`__fp_basics_pack_weird_low:NNNw`
`__fp_basics_pack_weird_high:NNNNNNNw`

This is used in `l3fp-basics` for additions and divisions. Their syntax is confusing, hence the name.

```

16532 \cs_new:Npn \__fp_basics_pack_weird_low:NNNw #1 #2#3#4 #5;
16533 {
16534     \if_meaning:w 2 #1
16535         + 1
16536     \fi:
16537     \__fp_int_eval_end:
16538     #2#3#4; {#5} ;
16539 }
16540 \cs_new:Npn \__fp_basics_pack_weird_high:NNNNNNNw
16541     1 #1#2#3#4 #5#6#7#8 #9; { ; {#1#2#3#4} {#5#6#7#8} {#9} }

```

(End definition for `__fp_basics_pack_weird_low:NNNw` and `__fp_basics_pack_weird_high:NNNNNNNw`.)

25.9 Decimate (dividing by a power of 10)

`__fp_decimate:nNnnnn`

`__fp_decimate:nNnnnn {⟨shift⟩} {f1}`
`{⟨X1⟩} {⟨X2⟩} {⟨X3⟩} {⟨X4⟩}`

Each $\langle X_i \rangle$ consists in 4 digits exactly, and $1000 \leq \langle X_1 \rangle < 9999$. The first argument determines by how much we shift the digits. $\langle f_1 \rangle$ is called as follows:

$\langle f_1 \rangle \langle \text{rounding} \rangle \{ \langle X'_1 \rangle \} \{ \langle X'_2 \rangle \} \langle \text{extra-digits} \rangle ;$

where $0 \leq \langle X'_i \rangle < 10^8 - 1$ are 8 digit integers, forming the truncation of our number. In other words,

$$\left(\sum_{i=1}^4 \langle X_i \rangle \cdot 10^{-4i} \cdot 10^{-\langle \text{shift} \rangle} \right) - (\langle X'_1 \rangle \cdot 10^{-8} + \langle X'_2 \rangle \cdot 10^{-16}) = 0. \langle \text{extra-digits} \rangle \cdot 10^{-16} \in [0, 10^{-16}).$$

To round properly later, we need to remember some information about the difference. The $\langle \text{rounding} \rangle$ digit is 0 if and only if the difference is exactly 0, and 5 if and only if the difference is exactly $0.5 \cdot 10^{-16}$. Otherwise, it is the (non-0, non-5) digit closest to 10^{17} times the difference. In particular, if the shift is 17 or more, all the digits are dropped, $\langle \text{rounding} \rangle$ is 1 (not 0), and $\langle X'_1 \rangle$ and $\langle X'_2 \rangle$ are both zero.

If the shift is 1, the $\langle \text{rounding} \rangle$ digit is simply the only digit that was pushed out of the brace groups (this is important for subtraction). It would be more natural for the $\langle \text{rounding} \rangle$ digit to be placed after the $\langle X'_i \rangle$, but the choice we make involves less reshuffling.

Note that this function treats negative $\langle \text{shift} \rangle$ as 0.

```

16542 \cs_new:Npn \__fp_decimate:nNnnnn #1
16543 {

```

```

16544 \cs:w
16545   __fp_decimate_
16546   \if_int_compare:w \__fp_int_eval:w #1 > \c__fp_prec_int
16547     tiny
16548   \else:
16549     \__fp_int_to_roman:w \__fp_int_eval:w #1
16550   \fi:
16551   :Nnnnn
16552 \cs_end:
16553 }

```

Each of the auxiliaries see the function $\langle f_1 \rangle$, followed by 4 blocks of 4 digits.

(End definition for `__fp_decimate:nNnnnn`.)

```

\__fp_decimate_:Nnnnn
\__fp_decimate_tiny:Nnnnn
16554 \cs_new:Npn \__fp_decimate_:Nnnnn #1 #2#3#4#5
16555   { #1 0 {#2#3} {#4#5} ; }
16556 \cs_new:Npn \__fp_decimate_tiny:Nnnnn #1 #2#3#4#5
16557   { #1 1 { 0000 0000 } { 0000 0000 } 0 #2#3#4#5 ; }

```

(End definition for `__fp_decimate_:Nnnnn` and `__fp_decimate_tiny:Nnnnn`.)

```

\__fp_decimate_auxi:Nnnnn
\__fp_decimate_auxii:Nnnnn
\__fp_decimate_auxiii:Nnnnn
\__fp_decimate_auxiv:Nnnnn
\__fp_decimate_auxv:Nnnnn
\__fp_decimate_auxvi:Nnnnn
\__fp_decimate_auxvii:Nnnnn
\__fp_decimate_auxviii:Nnnnn
\__fp_decimate_auxix:Nnnnn
\__fp_decimate_auxx:Nnnnn
\__fp_decimate_auxxi:Nnnnn
\__fp_decimate_auxxii:Nnnnn
\__fp_decimate_auxxiii:Nnnnn
\__fp_decimate_auxxiv:Nnnnn
\__fp_decimate_auxxv:Nnnnn
\__fp_decimate_auxxvi:Nnnnn

```

```

\__fp_decimate_auxi:Nnnnn <math>\langle f_1 \rangle</math> {<math>\langle X_1 \rangle</math>} {<math>\langle X_2 \rangle</math>} {<math>\langle X_3 \rangle</math>} {<math>\langle X_4 \rangle</math>}

```

Shifting happens in two steps: compute the $\langle \textit{rounding} \rangle$ digit, and repack digits into two blocks of 8. The sixteen functions are very similar, and defined through `__fp_tmp:w`. The arguments are as follows: #1 indicates which function is being defined; after one step of expansion, #2 yields the “extra digits” which are then converted by `__fp_round_digit:Nw` to the $\langle \textit{rounding} \rangle$ digit (note the + separating blocks of digits to avoid overflowing TeX’s integers). This triggers the f-expansion of `__fp_decimate_pack:nnnnnnnnnw`,⁹ responsible for building two blocks of 8 digits, and removing the rest. For this to work, #3 alternates between braced and unbraced blocks of 4 digits, in such a way that the 5 first and 5 next token groups yield the correct blocks of 8 digits.

```

16558 \cs_new:Npn \__fp_tmp:w #1 #2 #3
16559   {
16560     \cs_new:cpn { __fp_decimate_ #1 :Nnnnn } ##1 ##2##3##4##5
16561     {
16562       \exp_after:wN ##1
16563       \int_value:w
16564       \exp_after:wN \__fp_round_digit:Nw #2 ;
16565       \__fp_decimate_pack:nnnnnnnnnw #3 ;
16566     }
16567   }
16568 \__fp_tmp:w {i}   {\use_none:nnn   #50}{ 0{#2}#3{#4}#5      }
16569 \__fp_tmp:w {ii}  {\use_none:nn    #5 }{ 00{#2}#3{#4}#5      }
16570 \__fp_tmp:w {iii} {\use_none:n     #5 }{ 000{#2}#3{#4}#5     }
16571 \__fp_tmp:w {iv}  {                 #5 }{ {0000}#2{#3}#4 #5    }
16572 \__fp_tmp:w {v}   {\use_none:nnn   #4#5 }{ 0{0000}#2{#3}#4 #5    }
16573 \__fp_tmp:w {vi}  {\use_none:nn    #4#5 }{ 00{0000}#2{#3}#4 #5    }
16574 \__fp_tmp:w {vii} {\use_none:n     #4#5 }{ 000{0000}#2{#3}#4 #5    }
16575 \__fp_tmp:w {viii}{                 #4#5 }{ {0000}0000{#2}#3 #4 #5 }
16576 \__fp_tmp:w {ix}  {\use_none:nnn   #3#4+#5}{ 0{0000}0000{#2}#3 #4 #5 }

```

⁹No, the argument spec is not a mistake: the function calls an auxiliary to do half of the job.

```

16577 \__fp_tmp:w {x}    {\use_none:nn #3#4+#5}{ 00{0000}0000{#2}#3 #4 #5    }
16578 \__fp_tmp:w {xi}   {\use_none:n  #3#4+#5}{ 000{0000}0000{#2}#3 #4 #5    }
16579 \__fp_tmp:w {xii}  {\use_none:n  #3#4+#5}{ {0000}0000{0000}#2 #3 #4 #5    }
16580 \__fp_tmp:w {xiii}{\use_none:nnn#2#3+#4#5}{ 0{0000}0000{0000}#2 #3 #4 #5    }
16581 \__fp_tmp:w {xiv}   {\use_none:nn #2#3+#4#5}{ 00{0000}0000{0000}#2 #3 #4 #5    }
16582 \__fp_tmp:w {xv}    {\use_none:n  #2#3+#4#5}{ 000{0000}0000{0000}#2 #3 #4 #5    }
16583 \__fp_tmp:w {xvi}   {\use_none:n  #2#3+#4#5}{ {0000}0000{0000}0000 #2 #3 #4 #5 }

```

(End definition for __fp_decimate_auxi:Nnnnn and others.)

__fp_decimate_pack:nnnnnnnnnw

The computation of the *rounding* digit leaves an unfinished \int_value:w, which expands the following functions. This allows us to repack nicely the digits we keep. Those digits come as an alternation of unbraced and braced blocks of 4 digits, such that the first 5 groups of token consist in 4 single digits, and one brace group (in some order), and the next 5 have the same structure. This is followed by some digits and a semicolon.

```

16584 \cs_new:Npn \__fp_decimate_pack:nnnnnnnnnw #1#2#3#4#5
16585   { \__fp_decimate_pack:nnnnnw { #1#2#3#4#5 } }
16586 \cs_new:Npn \__fp_decimate_pack:nnnnnw #1 #2#3#4#5#6
16587   { {#1} {#2#3#4#5#6} }

```

(End definition for __fp_decimate_pack:nnnnnnnnnw.)

25.10 Functions for use within primitive conditional branches

The functions described in this section are not pretty and can easily be misused. When correctly used, each of them removes one \fi: as part of its parameter text, and puts one back as part of its replacement text.

Many computation functions in l3fp must perform tests on the type of floating points that they receive. This is often done in an \if_case:w statement or another conditional statement, and only a few cases lead to actual computations: most of the special cases are treated using a few standard functions which we define now. A typical use context for those functions would be

```

\if_case:w <integer> \exp_stop_f:
  \__fp_case_return_o:Nw <fp var>
\or: \__fp_case_use:nw {<some computation>}
\or: \__fp_case_return_same_o:w
\or: \__fp_case_return:nw {<something>}
\fi:
<junk>
<floating point>

```

In this example, the case 0 returns the floating point *<fp var>*, expanding once after that floating point. Case 1 does *<some computation>* using the *<floating point>* (presumably compute the operation requested by the user in that non-trivial case). Case 2 returns the *<floating point>* without modifying it, removing the *<junk>* and expanding once after. Case 3 closes the conditional, removes the *<junk>* and the *<floating point>*, and expands *<something>* next. In other cases, the “*<junk>*” is expanded, performing some other operation on the *<floating point>*. We provide similar functions with two trailing *<floating points>*.

`__fp_case_use:nw` This function ends a T_EX conditional, removes junk until the next floating point, and places its first argument before that floating point, to perform some operation on the floating point.

```
16588 \cs_new:Npn \__fp_case_use:nw #1#2 \fi: #3 \s__fp { \fi: #1 \s__fp }
```

(End definition for `__fp_case_use:nw`.)

`__fp_case_return:nw` This function ends a T_EX conditional, removes junk and a floating point, and places its first argument in the input stream. A quirk is that we don't define this function requiring a floating point to follow, simply anything ending in a semicolon. This, in turn, means that the *<junk>* may not contain semicolons.

```
16589 \cs_new:Npn \__fp_case_return:nw #1#2 \fi: #3 ; { \fi: #1 }
```

(End definition for `__fp_case_return:nw`.)

`__fp_case_return_o:Nw` This function ends a T_EX conditional, removes junk and a floating point, and returns its first argument (an *<fp var>*) then expands once after it.

```
16590 \cs_new:Npn \__fp_case_return_o:Nw #1#2 \fi: #3 \s__fp #4 ;
16591 { \fi: \exp_after:wN #1 }
```

(End definition for `__fp_case_return_o:Nw`.)

`__fp_case_return_same_o:w` This function ends a T_EX conditional, removes junk, and returns the following floating point, expanding once after it.

```
16592 \cs_new:Npn \__fp_case_return_same_o:w #1 \fi: #2 \s__fp
16593 { \fi: \__fp_exp_after_o:w \s__fp }
```

(End definition for `__fp_case_return_same_o:w`.)

`__fp_case_return_o:Nww` Same as `__fp_case_return_o:Nw` but with two trailing floating points.

```
16594 \cs_new:Npn \__fp_case_return_o:Nww #1#2 \fi: #3 \s__fp #4 ; #5 ;
16595 { \fi: \exp_after:wN #1 }
```

(End definition for `__fp_case_return_o:Nww`.)

`__fp_case_return_i_o:ww` Similar to `__fp_case_return_same_o:w`, but this returns the first or second of two trailing floating point numbers, expanding once after the result.

```
16596 \cs_new:Npn \__fp_case_return_i_o:ww #1 \fi: #2 \s__fp #3 ; \s__fp #4 ;
16597 { \fi: \__fp_exp_after_o:w \s__fp #3 ; }
16598 \cs_new:Npn \__fp_case_return_ii_o:ww #1 \fi: #2 \s__fp #3 ;
16599 { \fi: \__fp_exp_after_o:w }
```

(End definition for `__fp_case_return_i_o:ww` and `__fp_case_return_ii_o:ww`.)

25.11 Integer floating points

`_fp_int_p:w` Tests if the floating point argument is an integer. For normal floating point numbers, this holds if the rounding digit resulting from `_fp_decimate:nNnnnn` is 0.

`_fp_int:wTF`

```

16600 \prg_new_conditional:Npnn \_fp_int:w \s_fp \_fp_chk:w #1 #2 #3 #4;
16601 { TF , T , F , p }
16602 {
16603   \if_case:w #1 \exp_stop_f:
16604     \prg_return_true:
16605   \or:
16606     \if_charcode:w 0
16607       \_fp_decimate:nNnnnn { \c__fp_prec_int - #3 }
16608       \_fp_use_i_until_s:nw #4
16609       \prg_return_true:
16610     \else:
16611       \prg_return_false:
16612     \fi:
16613   \else: \prg_return_false:
16614   \fi:
16615 }
```

(End definition for `_fp_int:wTF`.)

25.12 Small integer floating points

`_fp_small_int:wTF` Tests if the floating point argument is an integer or $\pm\infty$. If so, it is clipped to an integer in the range $[-10^8, 10^8]$ and fed as a braced argument to the *⟨true code⟩*. Otherwise, the *⟨false code⟩* is performed.

`_fp_small_int_true:wTF`

`_fp_small_int_normal:NnwTF`

`_fp_small_int_test:NnnwNTF`

First filter special cases: zeros and infinities are integers, `nan` is not. For normal numbers, decimate. If the rounding digit is not 0 run the *⟨false code⟩*. If it is, then the integer is `#2 #3`; use `#3` if `#2` vanishes and otherwise 10^8 .

```

16616 \cs_new:Npn \_fp_small_int:wTF \s_fp \_fp_chk:w #1#2
16617 {
16618   \if_case:w #1 \exp_stop_f:
16619     \_fp_case_return:nw { \_fp_small_int_true:wTF 0 ; }
16620   \or: \exp_after:wN \_fp_small_int_normal:NnwTF
16621   \or:
16622     \_fp_case_return:nw
16623     {
16624       \exp_after:wN \_fp_small_int_true:wTF \int_value:w
16625       \if_meaning:w 2 #2 - \fi: 1 0000 0000 ;
16626     }
16627   \else: \_fp_case_return:nw \use_ii:nn
16628   \fi:
16629   #2
16630 }
16631 \cs_new:Npn \_fp_small_int_true:wTF #1; #2#3 { #2 {#1} }
16632 \cs_new:Npn \_fp_small_int_normal:NnwTF #1#2#3;
16633 {
16634   \_fp_decimate:nNnnnn { \c__fp_prec_int - #2 }
16635   \_fp_small_int_test:NnnwNw
16636   #3 #1
16637 }
```

```

16638 \cs_new:Npn \__fp_small_int_test:NnnwNw #1#2#3#4; #5
16639 {
16640   \if_meaning:w 0 #1
16641     \exp_after:wN \__fp_small_int_true:wTF
16642     \int_value:w \if_meaning:w 2 #5 - \fi:
16643     \if_int_compare:w #2 > 0 \exp_stop_f:
16644       1 0000 0000
16645     \else:
16646       #3
16647     \fi:
16648     \exp_after:wN ;
16649   \else:
16650     \exp_after:wN \use_ii:nn
16651   \fi:
16652 }

```

(End definition for __fp_small_int:wTF and others.)

25.13 Fast string comparison

__fp_str_if_eq:nn A private version of the low-level string comparison function.

```

16653 \cs_new_eq:NN \__fp_str_if_eq:nn \tex_strcmp:D

```

(End definition for __fp_str_if_eq:nn.)

25.14 Name of a function from its l3fp-parse name

__fp_func_to_name:N The goal is to convert for instance __fp_sin_o:w to sin. This is used in error messages
 __fp_func_to_name_aux:w hence does not need to be fast.

```

16654 \cs_new:Npn \__fp_func_to_name:N #1
16655 {
16656   \exp_last_unbraced:Nf
16657     \__fp_func_to_name_aux:w { \cs_to_str:N #1 } X
16658 }
16659 \cs_set_protected:Npn \__fp_tmp:w #1 #2
16660 { \cs_new:Npn \__fp_func_to_name_aux:w ##1 #1 ##2 #2 ##3 X {##2} }
16661 \exp_args:Nff \__fp_tmp:w { \tl_to_str:n { __fp_ } }
16662 { \tl_to_str:n { _o: } }

```

(End definition for __fp_func_to_name:N and __fp_func_to_name_aux:w.)

25.15 Messages

Using a floating point directly is an error.

```

16663 \__kernel_msg_new:nnnn { kernel } { misused-fp }
16664 { A~floating~point~with~value~'#1'~was~misused. }
16665 {
16666   To~obtain~the~value~of~a~floating~point~variable,~use~
16667   '\token_to_str:N \fp_to_decimal:N',~
16668   '\token_to_str:N \fp_to_tl:N',~or~other~
16669   conversion~functions.
16670 }
16671 </package>

```


26 13fp-traps Implementation

16672 `<*package>`

16673 `<@@=fp>`

Exceptions should be accessed by an `n`-type argument, among

- `invalid_operation`
- `division_by_zero`
- `overflow`
- `underflow`
- `inexact` (actually never used).

26.1 Flags

Flags to denote exceptions.

`flag_fp_invalid_operation`
`flag_fp_division_by_zero`
`flag_fp_overflow`
`flag_fp_underflow`

16674 `\flag_new:n { fp_invalid_operation }`

16675 `\flag_new:n { fp_division_by_zero }`

16676 `\flag_new:n { fp_overflow }`

16677 `\flag_new:n { fp_underflow }`

(End definition for flag `fp_invalid_operation` and others. These variables are documented on page 212.)

26.2 Traps

Exceptions can be trapped to obtain custom behaviour. When an invalid operation or a division by zero is trapped, the trap receives as arguments the result as an `N`-type floating point number, the function name (multiple letters for prefix operations, or a single symbol for infix operations), and the operand(s). When an overflow or underflow is trapped, the trap receives the resulting overly large or small floating point number if it is not too big, otherwise it receives $+\infty$. Currently, the `inexact` exception is entirely ignored.

The behaviour when an exception occurs is controlled by the definitions of the functions

- `__fp_invalid_operation:nnw,`
- `__fp_invalid_operation_o:Nww,`
- `__fp_invalid_operation_tl_o:ff,`
- `__fp_division_by_zero_o:Nnw,`
- `__fp_division_by_zero_o:NNww,`
- `__fp_overflow:w,`
- `__fp_underflow:w.`

Rather than changing them directly, we provide a user interface as `\fp_trap:nn {<exception>} {<way of trapping>}`, where the `<way of trapping>` is one of `error`, `flag`, or `none`.

We also provide `__fp_invalid_operation_o:nw`, defined in terms of `__fp_invalid_operation:nnw`.

`\fp_trap:nn`

```
16678 \cs_new_protected:Npn \fp_trap:nn #1#2
16679 {
16680   \cs_if_exist_use:cF { __fp_trap_#1_set_#2: }
16681   {
16682     \clist_if_in:nnTF
16683     { invalid_operation , division_by_zero , overflow , underflow }
16684     {#1}
16685     {
16686       \__kernel_msg_error:nnxx { kernel }
16687       { unknown-fpu-trap-type } {#1} {#2}
16688     }
16689     {
16690       \__kernel_msg_error:nnx
16691       { kernel } { unknown-fpu-exception } {#1}
16692     }
16693   }
16694 }
```

(End definition for `\fp_trap:nn`. This function is documented on page 212.)

`_fp_trap_invalid_operation_set_error:` We provide three types of trapping for invalid operations: either produce an error and
`_fp_trap_invalid_operation_set_flag:` raise the relevant flag; or only raise the flag; or don't even raise the flag. In most cases,
`_fp_trap_invalid_operation_set_none:` the function produces as a result its first argument, possibly with post-expansion.
`_fp_trap_invalid_operation_set:N`

```
16695 \cs_new_protected:Npn \_fp_trap_invalid_operation_set_error:
16696 { \_fp_trap_invalid_operation_set:N \prg_do_nothing: }
16697 \cs_new_protected:Npn \_fp_trap_invalid_operation_set_flag:
16698 { \_fp_trap_invalid_operation_set:N \use_none:nnnnn }
16699 \cs_new_protected:Npn \_fp_trap_invalid_operation_set_none:
16700 { \_fp_trap_invalid_operation_set:N \use_none:nnnnnnnn }
16701 \cs_new_protected:Npn \_fp_trap_invalid_operation_set:N #1
16702 {
16703   \exp_args:Nno \use:n
16704   { \cs_set:Npn \_fp_invalid_operation:nnw ##1##2##3; }
16705   {
16706     #1
16707     \_fp_error:nnfn { fp-invalid } {##2} { \fp_to_tl:n { ##3; } } { }
16708     \flag_raise_if_clear:n { fp_invalid_operation }
16709     ##1
16710   }
16711   \exp_args:Nno \use:n
16712   { \cs_set:Npn \_fp_invalid_operation_o:Nww ##1##2; ##3; }
16713   {
16714     #1
16715     \_fp_error:nffn { fp-invalid-ii }
16716     { \fp_to_tl:n { ##2; } } { \fp_to_tl:n { ##3; } } {##1}
16717     \flag_raise_if_clear:n { fp_invalid_operation }
16718     \exp_after:wN \c_nan_fp
16719   }
16720   \exp_args:Nno \use:n
16721   { \cs_set:Npn \_fp_invalid_operation_tl_o:ff ##1##2 }
16722   {
16723     #1
16724     \_fp_error:nffn { fp-invalid } {##1} {##2} { }
```

```

16725         \flag_raise_if_clear:n { fp_invalid_operation }
16726         \exp_after:wN \c_nan_fp
16727     }
16728 }

```

(End definition for `__fp_trap_invalid_operation_set_error:` and others.)

`__fp_trap_division_by_zero_set_error:` We provide three types of trapping for invalid operations and division by zero: either produce an error and raise the relevant flag; or only raise the flag; or don't even raise the flag. In all cases, the function must produce a result, namely its first argument, $\pm\infty$ or NaN.

```

16729 \cs_new_protected:Npn \__fp_trap_division_by_zero_set_error:
16730 { \__fp_trap_division_by_zero_set:N \prg_do_nothing: }
16731 \cs_new_protected:Npn \__fp_trap_division_by_zero_set_flag:
16732 { \__fp_trap_division_by_zero_set:N \use_none:nnnnn }
16733 \cs_new_protected:Npn \__fp_trap_division_by_zero_set_none:
16734 { \__fp_trap_division_by_zero_set:N \use_none:nnnnnnn }
16735 \cs_new_protected:Npn \__fp_trap_division_by_zero_set:N #1
16736 {
16737     \exp_args:Nno \use:n
16738     { \cs_set:Npn \__fp_division_by_zero_o:NNw ##1##2##3; }
16739     {
16740         #1
16741         \__fp_error:nfn { fp-zero-div } {##2} { \fp_to_tl:n { ##3; } } { }
16742         \flag_raise_if_clear:n { fp_division_by_zero }
16743         \exp_after:wN ##1
16744     }
16745     \exp_args:Nno \use:n
16746     { \cs_set:Npn \__fp_division_by_zero_o:NNww ##1##2##3; ##4; }
16747     {
16748         #1
16749         \__fp_error:nfn { fp-zero-div-ii }
16750         { \fp_to_tl:n { ##3; } } { \fp_to_tl:n { ##4; } } {##2}
16751         \flag_raise_if_clear:n { fp_division_by_zero }
16752         \exp_after:wN ##1
16753     }
16754 }

```

(End definition for `__fp_trap_division_by_zero_set_error:` and others.)

`__fp_trap_overflow_set_error:` Just as for invalid operations and division by zero, the three different behaviours are obtained by feeding `\prg_do_nothing:`, `\use_none:nnnnn` or `\use_none:nnnnnnn` to an auxiliary, with a further auxiliary common to overflow and underflow functions. In most cases, the argument of the `__fp_overflow:w` and `__fp_underflow:w` functions will be an (almost) normal number (with an exponent outside the allowed range), and the error message thus displays that number together with the result to which it overflowed or underflowed. For extreme cases such as `10 ** 1e9999`, the exponent would be too large for T_EX, and `__fp_overflow:w` receives $\pm\infty$ (`__fp_underflow:w` would receive ± 0); then we cannot do better than simply say an overflow or underflow occurred.

```

16755 \cs_new_protected:Npn \__fp_trap_overflow_set_error:
16756 { \__fp_trap_overflow_set:N \prg_do_nothing: }
16757 \cs_new_protected:Npn \__fp_trap_overflow_set_flag:
16758 { \__fp_trap_overflow_set:N \use_none:nnnnn }
16759 \cs_new_protected:Npn \__fp_trap_overflow_set_none:

```

```

16760 { \__fp_trap_overflow_set:N \use_none:nnnnnnn }
16761 \cs_new_protected:Npn \__fp_trap_overflow_set:N #1
16762 { \__fp_trap_overflow_set:NnNn #1 { overflow } \__fp_inf_fp:N { inf } }
16763 \cs_new_protected:Npn \__fp_trap_underflow_set_error:
16764 { \__fp_trap_underflow_set:N \prg_do_nothing: }
16765 \cs_new_protected:Npn \__fp_trap_underflow_set_flag:
16766 { \__fp_trap_underflow_set:N \use_none:nnnnn }
16767 \cs_new_protected:Npn \__fp_trap_underflow_set_none:
16768 { \__fp_trap_underflow_set:N \use_none:nnnnnnn }
16769 \cs_new_protected:Npn \__fp_trap_underflow_set:N #1
16770 { \__fp_trap_overflow_set:NnNn #1 { underflow } \__fp_zero_fp:N { 0 } }
16771 \cs_new_protected:Npn \__fp_trap_overflow_set:NnNn #1#2#3#4
16772 {
16773   \exp_args:Nno \use:n
16774   { \cs_set:cpn { \__fp_ #2 :w } \s__fp \__fp_chk:w ##1##2##3; }
16775   {
16776     #1
16777     \__fp_error:nffn
16778     { fp-flow \if_meaning:w 1 ##1 -to \fi: }
16779     { \fp_to_tl:n { \s__fp \__fp_chk:w ##1##2##3; } }
16780     { \token_if_eq_meaning:NNF 0 ##2 { - } #4 }
16781     {#2}
16782     \flag_raise_if_clear:n { fp_#2 }
16783     #3 ##2
16784   }
16785 }

```

(End definition for __fp_trap_overflow_set_error: and others.)

__fp_invalid_operation:nnw Initialize the control sequences (to log properly their existence). Then set invalid operations to trigger an error, and division by zero, overflow, and underflow to act silently on their flag.

```

\__fp_invalid_operation_o:Nnw
\__fp_invalid_operation_tl_o:ff
\__fp_division_by_zero_o:Nnw
\__fp_division_by_zero_o:NNww
\__fp_overflow:w
\__fp_underflow:w
16786 \cs_new:Npn \__fp_invalid_operation:nnw #1#2#3; { }
16787 \cs_new:Npn \__fp_invalid_operation_o:Nnw #1#2; #3; { }
16788 \cs_new:Npn \__fp_invalid_operation_tl_o:ff #1 #2 { }
16789 \cs_new:Npn \__fp_division_by_zero_o:Nnw #1#2#3; { }
16790 \cs_new:Npn \__fp_division_by_zero_o:NNww #1#2#3; #4; { }
16791 \cs_new:Npn \__fp_overflow:w { }
16792 \cs_new:Npn \__fp_underflow:w { }
16793 \fp_trap:nn { invalid_operation } { error }
16794 \fp_trap:nn { division_by_zero } { flag }
16795 \fp_trap:nn { overflow } { flag }
16796 \fp_trap:nn { underflow } { flag }

```

(End definition for __fp_invalid_operation:nnw and others.)

__fp_invalid_operation_o:nw Convenient short-hands for returning \c_nan_fp for a unary or binary operation, and expanding after.

```

\__fp_invalid_operation_o:nw
\__fp_invalid_operation_o:fw
16797 \cs_new:Npn \__fp_invalid_operation_o:nw
16798 { \__fp_invalid_operation:nnw { \exp_after:wN \c_nan_fp } }
16799 \cs_generate_variant:Nn \__fp_invalid_operation_o:nw { f }

```

(End definition for __fp_invalid_operation_o:nw.)

26.3 Errors

```

\__fp_error:nnnn
\__fp_error:nnfn 16800 \cs_new:Npn \__fp_error:nnnn
\__fp_error:nffn 16801 { \__kernel_msg_expandable_error:nnnnn { kernel } }
\__fp_error:nfff 16802 \cs_generate_variant:Nn \__fp_error:nnnn { nnf, nff , nfff }

(End definition for \__fp_error:nnnn.)

```

26.4 Messages

Some messages.

```

16803 \__kernel_msg_new:nnnn { kernel } { unknown-fpu-exception }
16804 {
16805   The~FPU~exception~'#1'~is~not~known:~
16806   that~trap~will~never~be~triggered.
16807 }
16808 {
16809   The~only~exceptions~to~which~traps~can~be~attached~are \
16810   \iow_indent:n
16811   {
16812     * ~ invalid_operation \
16813     * ~ division_by_zero \
16814     * ~ overflow \
16815     * ~ underflow
16816   }
16817 }
16818 \__kernel_msg_new:nnnn { kernel } { unknown-fpu-trap-type }
16819 { The~FPU~trap~type~'#2'~is~not~known. }
16820 {
16821   The~trap~type~must~be~one~of \
16822   \iow_indent:n
16823   {
16824     * ~ error \
16825     * ~ flag \
16826     * ~ none
16827   }
16828 }
16829 \__kernel_msg_new:nnn { kernel } { fp-flow }
16830 { An ~ #3 ~ occurred. }
16831 \__kernel_msg_new:nnn { kernel } { fp-flow-to }
16832 { #1 ~ #3 ed ~ to ~ #2 . }
16833 \__kernel_msg_new:nnn { kernel } { fp-zero-div }
16834 { Division~by~zero~in~ #1 (#2) }
16835 \__kernel_msg_new:nnn { kernel } { fp-zero-div-ii }
16836 { Division~by~zero~in~ (#1) #3 (#2) }
16837 \__kernel_msg_new:nnn { kernel } { fp-invalid }
16838 { Invalid~operation~ #1 (#2) }
16839 \__kernel_msg_new:nnn { kernel } { fp-invalid-ii }
16840 { Invalid~operation~ (#1) #3 (#2) }
16841 \__kernel_msg_new:nnn { kernel } { fp-unknown-type }
16842 { Unknown~type~for~'#1' }
16843 \end{package}

```

27 13fp-round implementation

```

16844 (*package)
16845 (@@=fp)

__fp_parse_word_trunc:N
__fp_parse_word_floor:N
__fp_parse_word_ceil:N
16846 \cs_new:Npn __fp_parse_word_trunc:N
16847 { __fp_parse_function:NNN __fp_round_o:Nw __fp_round_to_zero:NNN }
16848 \cs_new:Npn __fp_parse_word_floor:N
16849 { __fp_parse_function:NNN __fp_round_o:Nw __fp_round_to_ninf:NNN }
16850 \cs_new:Npn __fp_parse_word_ceil:N
16851 { __fp_parse_function:NNN __fp_round_o:Nw __fp_round_to_pinf:NNN }

(End definition for __fp_parse_word_trunc:N, __fp_parse_word_floor:N, and __fp_parse_word_ceil:N.)

```

```

__fp_parse_word_round:N
__fp_parse_round:Nw
16852 \cs_new:Npn __fp_parse_word_round:N #1#2
16853 {
16854   __fp_parse_function:NNN
16855   __fp_round_o:Nw __fp_round_to_nearest:NNN #1
16856   #2
16857 }
16858 \cs_new:Npn __fp_parse_round:Nw #1 #2 __fp_round_to_nearest:NNN #3#4
16859 { #2 #1 #3 }
16860

(End definition for __fp_parse_word_round:N and __fp_parse_round:Nw.)

```

27.1 Rounding tools

`\c__fp_five_int` This is used as the half-point for which numbers are rounded up/down.

```

16861 \int_const:Nn \c__fp_five_int { 5 }

```

(End definition for `\c__fp_five_int`.)

Floating point operations often yield a result that cannot be exactly represented in a significand with 16 digits. In that case, we need to round the exact result to a representable number. The IEEE standard defines four rounding modes:

- Round to nearest: round to the representable floating point number whose absolute difference with the exact result is the smallest. If the exact result lies exactly at the mid-point between two consecutive representable floating point numbers, round to the floating point number whose last digit is even.
- Round towards negative infinity: round to the greatest floating point number not larger than the exact result.
- Round towards zero: round to a floating point number with the same sign as the exact result, with the largest absolute value not larger than the absolute value of the exact result.
- Round towards positive infinity: round to the least floating point number not smaller than the exact result.

This is not fully implemented in l3fp yet, and transcendental functions fall back on the “round to nearest” mode. All rounding for basic algebra is done through the functions defined in this module, which can be redefined to change their rounding behaviour (but there is not interface for that yet).

The rounding tools available in this module are many variations on a base function `__fp_round:NNN`, which expands to `0\exp_stop_f:` or `1\exp_stop_f:` depending on whether the final result should be rounded up or down.

- `__fp_round:NNN <sign> <digit1> <digit2>` can expand to `0\exp_stop_f:` or `1\exp_stop_f:.`
- `__fp_round_s:NNNw <sign> <digit1> <digit2> <more digits>`; can expand to `0\exp_stop_f:;` or `1\exp_stop_f:;`.
- `__fp_round_neg:NNN <sign> <digit1> <digit2>` can expand to `0\exp_stop_f:` or `1\exp_stop_f:.`

See implementation comments for details on the syntax.

```
\__fp_round:NNN
\__fp_round_to_nearest:NNN
  \__fp_round_to_nearest_ninf:NNN
  \__fp_round_to_nearest_zero:NNN
  \__fp_round_to_nearest_pinf:NNN
\__fp_round_to_ninf:NNN
\__fp_round_to_zero:NNN
\__fp_round_to_pinf:NNN
```

`__fp_round:NNN <final sign> <digit1> <digit2>`
If rounding the number $\langle final\ sign \rangle \langle digit_1 \rangle . \langle digit_2 \rangle$ to an integer rounds it towards zero (truncates it), this function expands to `0\exp_stop_f:`, and otherwise to `1\exp_stop_f:.` Typically used within the scope of an `__fp_int_eval:w`, to add 1 if needed, and thereby round correctly. The result depends on the rounding mode.

It is very important that $\langle final\ sign \rangle$ be the final sign of the result. Otherwise, the result would be incorrect in the case of rounding towards $-\infty$ or towards $+\infty$. Also recall that $\langle final\ sign \rangle$ is 0 for positive, and 2 for negative.

By default, the functions below return `0\exp_stop_f:`, but this is superseded by `__fp_round_return_one:`, which instead returns `1\exp_stop_f:`, expanding everything and removing `0\exp_stop_f:` in the process. In the case of rounding towards $\pm\infty$ or towards 0, this is not really useful, but it prepares us for the “round to nearest, ties to even” mode.

The “round to nearest” mode is the default. If the $\langle digit_2 \rangle$ is larger than 5, then round up. If it is less than 5, round down. If it is exactly 5, then round such that $\langle digit_1 \rangle$ plus the result is even. In other words, round up if $\langle digit_1 \rangle$ is odd.

The “round to nearest” mode has three variants, which differ in how ties are rounded: down towards $-\infty$, truncated towards 0, or up towards $+\infty$.

```
16862 \cs_new:Npn \__fp_round_return_one:
16863   { \exp_after:wN 1 \exp_after:wN \exp_stop_f: \exp:w }
16864 \cs_new:Npn \__fp_round_to_ninf:NNN #1 #2 #3
16865   {
16866     \if_meaning:w 2 #1
16867       \if_int_compare:w #3 > 0 \exp_stop_f:
16868         \__fp_round_return_one:
16869       \fi:
16870     \fi:
16871     0 \exp_stop_f:
16872   }
16873 \cs_new:Npn \__fp_round_to_zero:NNN #1 #2 #3 { 0 \exp_stop_f: }
16874 \cs_new:Npn \__fp_round_to_pinf:NNN #1 #2 #3
16875   {
16876     \if_meaning:w 0 #1
16877       \if_int_compare:w #3 > 0 \exp_stop_f:
```

```

16878         \_fp_round_return_one:
16879         \fi:
16880         \fi:
16881         0 \exp_stop_f:
16882     }
16883 \cs_new:Npn \_fp_round_to_nearest:NNN #1 #2 #3
16884 {
16885     \if_int_compare:w #3 > \c__fp_five_int
16886         \_fp_round_return_one:
16887     \else:
16888         \if_meaning:w 5 #3
16889             \if_int_odd:w #2 \exp_stop_f:
16890             \_fp_round_return_one:
16891         \fi:
16892     \fi:
16893     \fi:
16894     0 \exp_stop_f:
16895 }
16896 \cs_new:Npn \_fp_round_to_nearest_ninf:NNN #1 #2 #3
16897 {
16898     \if_int_compare:w #3 > \c__fp_five_int
16899         \_fp_round_return_one:
16900     \else:
16901         \if_meaning:w 5 #3
16902             \if_meaning:w 2 #1
16903                 \_fp_round_return_one:
16904             \fi:
16905         \fi:
16906     \fi:
16907     0 \exp_stop_f:
16908 }
16909 \cs_new:Npn \_fp_round_to_nearest_zero:NNN #1 #2 #3
16910 {
16911     \if_int_compare:w #3 > \c__fp_five_int
16912         \_fp_round_return_one:
16913     \fi:
16914     0 \exp_stop_f:
16915 }
16916 \cs_new:Npn \_fp_round_to_nearest_pinf:NNN #1 #2 #3
16917 {
16918     \if_int_compare:w #3 > \c__fp_five_int
16919         \_fp_round_return_one:
16920     \else:
16921         \if_meaning:w 5 #3
16922             \if_meaning:w 0 #1
16923                 \_fp_round_return_one:
16924             \fi:
16925         \fi:
16926     \fi:
16927     0 \exp_stop_f:
16928 }
16929 \cs_new_eq:NN \_fp_round:NNN \_fp_round_to_nearest:NNN

```

(End definition for _fp_round:NNN and others.)

_fp_round_s:NNNw

_fp_round_s:NNNw <final sign> <digit> <more digits> ;

Similar to _fp_round:NNN, but with an extra semicolon, this function expands to 0\exp_stop_f:; if rounding <final sign><digit>.<more digits> to an integer truncates, and to 1\exp_stop_f:; otherwise. The <more digits> part must be a digit, followed by something that does not overflow a \int_use:N _fp_int_eval:w construction. The only relevant information about this piece is whether it is zero or not.

```

16930 \cs_new:Npn \_fp\_round\_s:NNNw #1 #2 #3 #4;
16931 {
16932   \exp\_after:wN \_fp\_round:NNN
16933   \exp\_after:wN #1
16934   \exp\_after:wN #2
16935   \int\_value:w \_fp\_int\_eval:w
16936   \if\_int\_odd:w 0 \if\_meaning:w 0 #3 1 \fi:
16937   \if\_meaning:w 5 #3 1 \fi:
16938   \exp\_stop\_f:
16939   \if\_int\_compare:w \_fp\_int\_eval:w #4 > 0 \exp\_stop\_f:
16940   1 +
16941   \fi:
16942   \fi:
16943   #3
16944   ;
16945 }

```

(End definition for _fp_round_s:NNNw.)

_fp_round_digit:Nw

\int_value:w _fp_round_digit:Nw <digit> <intexpr> ;

This function should always be called within an \int_value:w or _fp_int_eval:w expansion; it may add an extra _fp_int_eval:w, which means that the integer or integer expression should not be ended with a synonym of \relax, but with a semi-colon for instance.

```

16946 \cs_new:Npn \_fp\_round\_digit:Nw #1 #2;
16947 {
16948   \if\_int\_odd:w \if\_meaning:w 0 #1 1 \else:
16949   \if\_meaning:w 5 #1 1 \else:
16950   0 \fi: \fi: \exp\_stop\_f:
16951   \if\_int\_compare:w \_fp\_int\_eval:w #2 > 0 \exp\_stop\_f:
16952   \_fp\_int\_eval:w 1 +
16953   \fi:
16954   \fi:
16955   #1
16956 }

```

(End definition for _fp_round_digit:Nw.)

_fp_round_neg:NNN

_fp_round_neg:NNN <final sign> <digit₁> <digit₂>

_fp_round_to_nearest_neg:NNN

This expands to 0\exp_stop_f: or 1\exp_stop_f: after doing the following test.

_fp_round_to_nearest_ninf_neg:NNN

Starting from a number of the form <final sign>0.<15 digits><digit₁> with exactly 15 (non-all-zero) digits before <digit₁>, subtract from it <final sign>0.0...0<digit₂>, where there are 16 zeros. If in the current rounding mode the result should be rounded down, then this function returns 1\exp_stop_f:. Otherwise, i.e., if the result is rounded back to the first operand, then this function returns 0\exp_stop_f:.

_fp_round_to_nearest_zero_neg:NNN

_fp_round_to_nearest_pinf_neg:NNN

_fp_round_to_ninf_neg:NNN

_fp_round_to_zero_neg:NNN

_fp_round_to_pinf_neg:NNN

It turns out that this negative “round to nearest” is identical to the positive one. And this is the default mode.

```

16957 \cs_new_eq:NN \__fp_round_to_ninf_neg:NNN \__fp_round_to_pinf:NNN
16958 \cs_new:Npn \__fp_round_to_zero_neg:NNN #1 #2 #3
16959 {
16960     \if_int_compare:w #3 > 0 \exp_stop_f:
16961     \__fp_round_return_one:
16962     \fi:
16963     0 \exp_stop_f:
16964 }
16965 \cs_new_eq:NN \__fp_round_to_pinf_neg:NNN \__fp_round_to_ninf:NNN
16966 \cs_new_eq:NN \__fp_round_to_nearest_neg:NNN \__fp_round_to_nearest:NNN
16967 \cs_new_eq:NN \__fp_round_to_nearest_ninf_neg:NNN
16968 \__fp_round_to_nearest_pinf:NNN
16969 \cs_new:Npn \__fp_round_to_nearest_zero_neg:NNN #1 #2 #3
16970 {
16971     \if_int_compare:w #3 < \c__fp_five_int \else:
16972     \__fp_round_return_one:
16973     \fi:
16974     0 \exp_stop_f:
16975 }
16976 \cs_new_eq:NN \__fp_round_to_nearest_pinf_neg:NNN
16977 \__fp_round_to_nearest_ninf:NNN
16978 \cs_new_eq:NN \__fp_round_neg:NNN \__fp_round_to_nearest_neg:NNN

```

(End definition for __fp_round_neg:NNN and others.)

27.2 The round function

__fp_round_o:Nw
__fp_round_aux_o:Nw

First check that all arguments are floating point numbers. The `trunc`, `ceil` and `floor` functions expect one or two arguments (the second is 0 by default), and the `round` function also accepts a third argument (`nan` by default), which changes `#1` from `__fp_round_to_nearest:NNN` to one of its analogues.

```

16979 \cs_new:Npn \__fp_round_o:Nw #1
16980 {
16981     \__fp_parse_function_all_fp_o:fnw
16982     { \__fp_round_name_from_cs:N #1 }
16983     { \__fp_round_aux_o:Nw #1 }
16984 }
16985 \cs_new:Npn \__fp_round_aux_o:Nw #1#2 @
16986 {
16987     \if_case:w
16988     \__fp_int_eval:w \__fp_array_count:n {#2} \__fp_int_eval_end:
16989     \__fp_round_no_arg_o:Nw #1 \exp:w
16990     \or: \__fp_round:Nwn #1 #2 {0} \exp:w
16991     \or: \__fp_round:Nww #1 #2 \exp:w
16992     \else: \__fp_round:Nwww #1 #2 @ \exp:w
16993     \fi:
16994     \exp_after:wN \exp_end:
16995 }

```

(End definition for __fp_round_o:Nw and __fp_round_aux_o:Nw.)

__fp_round_no_arg_o:Nw

```

16996 \cs_new:Npn \__fp_round_no_arg_o:Nw #1
16997 {

```

```

16998 \cs_if_eq:NNTF #1 \__fp_round_to_nearest:NNN
16999 { \__fp_error:nnnn { fp-num-args } { round () } { 1 } { 3 } }
17000 {
17001   \__fp_error:nffn { fp-num-args }
17002   { \__fp_round_name_from_cs:N #1 () } { 1 } { 2 }
17003 }
17004 \exp_after:wN \c_nan_fp
17005 }

```

(End definition for __fp_round_no_arg_o:Nw.)

__fp_round:Nwww Having three arguments is only allowed for round, not trunc, ceil, floor, so check for that case. If all is well, construct one of __fp_round_to_nearest:NNN, __fp_round_to_nearest_zero:NNN, __fp_round_to_nearest_ninf:NNN, __fp_round_to_nearest_pinf:NNN and act accordingly.

```

17006 \cs_new:Npn \__fp_round:Nwww #1#2 ; #3 ; \s__fp \__fp_chk:w #4#5#6 ; #7 @
17007 {
17008   \cs_if_eq:NNTF #1 \__fp_round_to_nearest:NNN
17009   {
17010     \tl_if_empty:nTF {#7}
17011     {
17012       \exp_args:Nc \__fp_round:Nww
17013       {
17014         __fp_round_to_nearest
17015         \if_meaning:w 0 #4 _zero \else:
17016         \if_case:w #5 \exp_stop_f: _pinf \or: \else: _ninf \fi: \fi:
17017         :NNN
17018       }
17019       #2 ; #3 ;
17020     }
17021     {
17022       \__fp_error:nnnn { fp-num-args } { round () } { 1 } { 3 }
17023       \exp_after:wN \c_nan_fp
17024     }
17025   }
17026   {
17027     \__fp_error:nffn { fp-num-args }
17028     { \__fp_round_name_from_cs:N #1 () } { 1 } { 2 }
17029     \exp_after:wN \c_nan_fp
17030   }
17031 }

```

(End definition for __fp_round:Nwww.)

__fp_round_name_from_cs:N

```

17032 \cs_new:Npn \__fp_round_name_from_cs:N #1
17033 {
17034   \cs_if_eq:NNTF #1 \__fp_round_to_zero:NNN { trunc }
17035   {
17036     \cs_if_eq:NNTF #1 \__fp_round_to_ninf:NNN { floor }
17037     {
17038       \cs_if_eq:NNTF #1 \__fp_round_to_pinf:NNN { ceil }
17039       { round }
17040     }
17041   }

```

```

17041     }
17042 }

```

(End definition for _fp_round_name_from_cs:N.)

_fp_round:Nww

_fp_round:Nwn

If the number of digits to round to is an integer or infinity all is good; if it is nan then just produce a nan; otherwise invalid as we have something like round(1,3.14) where the number of digits is not an integer.

_fp_round_normal:NwNNnw

_fp_round_normal:NnnwNNnn

```

17043 \cs_new:Npn \_fp\_round:Nww #1#2 ; #3 ;

```

_fp_round_pack:Nw

```

17044 {

```

_fp_round_normal:NNwNnn

```

17045   \_fp\_small\_int:wTF #3; { \_fp\_round:Nwn #1#2; }

```

_fp_round_normal_end:wwNnn

```

17046   {

```

_fp_round_special:NwwNnn

```

17047     \if:w 3 \_fp\_kind:w #3 ;

```

_fp_round_special_aux:Nw

```

17048     \exp\_after:wN \use\_i:nn

```

```

17049     \else:

```

```

17050       \exp\_after:wN \use\_ii:nn

```

```

17051       \fi:

```

```

17052       { \exp\_after:wN \c\_nan\_fp }

```

```

17053       {

```

```

17054         \_fp\_invalid\_operation\_tl\_o:ff

```

```

17055         { \_fp\_round\_name\_from\_cs:N #1 }

```

```

17056         { \_fp\_array\_to\_clist:n { #2; #3; } }

```

```

17057       }

```

```

17058     }

```

```

17059   }

```

```

17060 \cs_new:Npn \_fp\_round:Nwn #1 \s\_fp \_fp\_chk:w #2#3#4; #5

```

```

17061 {

```

```

17062   \if\_meaning:w 1 #2

```

```

17063     \exp\_after:wN \_fp\_round\_normal:NwNNnw

```

```

17064     \exp\_after:wN #1

```

```

17065     \int\_value:w #5

```

```

17066   \else:

```

```

17067     \exp\_after:wN \_fp\_exp\_after\_o:w

```

```

17068     \fi:

```

```

17069     \s\_fp \_fp\_chk:w #2#3#4;

```

```

17070   }

```

```

17071 \cs_new:Npn \_fp\_round\_normal:NwNNnw #1#2 \s\_fp \_fp\_chk:w 1#3#4#5;

```

```

17072 {

```

```

17073   \_fp\_decimate:nNnnnn { \c\_fp\_prec\_int - #4 - #2 }

```

```

17074   \_fp\_round\_normal:NnnwNNnn #5 #1 #3 {#4} {#2}

```

```

17075 }

```

```

17076 \cs_new:Npn \_fp\_round\_normal:NnnwNNnn #1#2#3#4; #5#6

```

```

17077 {

```

```

17078   \exp\_after:wN \_fp\_round\_normal:NNwNnn

```

```

17079   \int\_value:w \_fp\_int\_eval:w

```

```

17080   \if\_int\_compare:w #2 > 0 \exp\_stop\_f:

```

```

17081     1 \int\_value:w #2

```

```

17082     \exp\_after:wN \_fp\_round\_pack:Nw

```

```

17083     \int\_value:w \_fp\_int\_eval:w 1#3 +

```

```

17084   \else:

```

```

17085     \if\_int\_compare:w #3 > 0 \exp\_stop\_f:

```

```

17086       1 \int\_value:w #3 +

```

```

17087     \fi:

```

```

17088   \fi:

```

```

17089     \exp_after:wN #5
17090     \exp_after:wN #6
17091     \use_none:nnnnnn #3
17092     #1
17093     \__fp_int_eval_end:
17094     0000 0000 0000 0000 ; #6
17095 }
17096 \cs_new:Npn \__fp_round_pack:Nw #1
17097 { \if_meaning:w 2 #1 + 1 \fi: \__fp_int_eval_end: }
17098 \cs_new:Npn \__fp_round_normal:NNwNnn #1 #2
17099 {
17100     \if_meaning:w 0 #2
17101     \exp_after:wN \__fp_round_special:NwwNnn
17102     \exp_after:wN #1
17103     \fi:
17104     \__fp_pack_twice_four:wNNNNNNNNN
17105     \__fp_pack_twice_four:wNNNNNNNNN
17106     \__fp_round_normal_end:wwNnn
17107     ; #2
17108 }
17109 \cs_new:Npn \__fp_round_normal_end:wwNnn #1;#2;#3#4#5
17110 {
17111     \exp_after:wN \__fp_exp_after_o:w \exp:w \exp_end_continue_f:w
17112     \__fp_sanitizew:Nw #3 #4 ; #1 ;
17113 }
17114 \cs_new:Npn \__fp_round_special:NwwNnn #1#2;#3;#4#5#6
17115 {
17116     \if_meaning:w 0 #1
17117     \__fp_case_return:nw
17118     { \exp_after:wN \__fp_zero_fp:N \exp_after:wN #4 }
17119     \else:
17120     \exp_after:wN \__fp_round_special_aux:Nw
17121     \exp_after:wN #4
17122     \int_value:w \__fp_int_eval:w 1
17123     \if_meaning:w 1 #1 -#6 \else: +#5 \fi:
17124     \fi:
17125     ;
17126 }
17127 \cs_new:Npn \__fp_round_special_aux:Nw #1#2;
17128 {
17129     \exp_after:wN \__fp_exp_after_o:w \exp:w \exp_end_continue_f:w
17130     \__fp_sanitizew:Nw #1#2; {1000}{0000}{0000}{0000};
17131 }

```

(End definition for __fp_round:Nww and others.)

```

17132 </package>

```

28 l3fp-parse implementation

```

17133 <*package>
17134 <@@=fp>

```

28.1 Work plan

The task at hand is non-trivial, and some previous failed attempts show that the code leads to unreadable logs, so we had better get it (almost) right the first time. Let us first describe our goal, then discuss the design precisely before writing any code.

In this file at least, a *floating point object* is a floating point number or tuple. This can be extended to anything that starts with `\s__fp` or `\s__fp_⟨type⟩` and ends with `;` with some internal structure that depends on the *type*.

`__fp_parse:n`

`__fp_parse:n {⟨fexpr⟩}`

Evaluates the *floating point expression* and leaves the result in the input stream as a floating point object. This function forms the basis of almost all public `l3fp` functions. During evaluation, each token is fully `f`-expanded.

`__fp_parse_o:n` does the same but expands once after its result.

T_EXhackers note: Registers (integers, toks, etc.) are automatically unpacked, without requiring a function such as `\int_use:N`. Invalid tokens remaining after `f`-expansion lead to unrecoverable low-level T_EX errors.

(End definition for `__fp_parse:n`.)

`\c__fp_prec_func_int`
`\c__fp_prec_hatii_int`
`\c__fp_prec_hat_int`
`\c__fp_prec_not_int`
`\c__fp_prec_juxt_int`
`\c__fp_prec_times_int`
`\c__fp_prec_plus_int`
`\c__fp_prec_comp_int`
`\c__fp_prec_and_int`
`\c__fp_prec_or_int`
`\c__fp_prec_quest_int`
`\c__fp_prec_colon_int`
`\c__fp_prec_comma_int`
`\c__fp_prec_tuple_int`
`\c__fp_prec_end_int`

Floating point expressions are composed of numbers, given in various forms, infix operators, such as `+`, `**`, or `,` (which joins two numbers into a list), and prefix operators, such as the unary `-`, functions, or opening parentheses. Here is a list of precedences which control the order of evaluation (some distinctions are irrelevant for the order of evaluation, but serve as signals), from the tightest binding to the loosest binding.

16 Function calls.

13/14 Binary `**` and `^` (right to left).

12 Unary `+`, `-`, `!` (right to left).

11 Juxtaposition (implicit `*`) with no parenthesis.

10 Binary `*` and `/`.

9 Binary `+` and `-`.

7 Comparisons.

6 Logical `and`, denoted by `&&`.

5 Logical `or`, denoted by `||`.

4 Ternary operator `?:`, piece `?`.

3 Ternary operator `?:`, piece `:`.

2 Commas.

1 Place where a comma is allowed and generates a tuple.

0 Start and end of the expression.

```

17135 \int_const:Nn \c__fp_prec_func_int { 16 }
17136 \int_const:Nn \c__fp_prec_hatii_int { 14 }
17137 \int_const:Nn \c__fp_prec_hat_int { 13 }
17138 \int_const:Nn \c__fp_prec_not_int { 12 }
17139 \int_const:Nn \c__fp_prec_juxt_int { 11 }
17140 \int_const:Nn \c__fp_prec_times_int { 10 }
17141 \int_const:Nn \c__fp_prec_plus_int { 9 }
17142 \int_const:Nn \c__fp_prec_comp_int { 7 }
17143 \int_const:Nn \c__fp_prec_and_int { 6 }
17144 \int_const:Nn \c__fp_prec_or_int { 5 }
17145 \int_const:Nn \c__fp_prec_quest_int { 4 }
17146 \int_const:Nn \c__fp_prec_colon_int { 3 }
17147 \int_const:Nn \c__fp_prec_comma_int { 2 }
17148 \int_const:Nn \c__fp_prec_tuple_int { 1 }
17149 \int_const:Nn \c__fp_prec_end_int { 0 }

```

(End definition for `\c__fp_prec_func_int` and others.)

28.1.1 Storing results

The main question in parsing expressions expandably is to decide where to put the intermediate results computed for various subexpressions.

One option is to store the values at the start of the expression, and carry them together as the first argument of each macro. However, we want to `f-expand` tokens one by one in the expression (as `\int_eval:n` does), and with this approach, expanding the next unread token forces us to jump with `\exp_after:wN` over every value computed earlier in the expression. With this approach, the run-time grows at least quadratically in the length of the expression, if not as its cube (inserting the `\exp_after:wN` is tricky and slow).

A second option is to place those values at the end of the expression. Then expanding the next unread token is straightforward, but this still hits a performance issue: for long expressions we would be reaching all the way to the end of the expression at every step of the calculation. The run-time is again quadratic.

A variation of the above attempts to place the intermediate results which appear when computing a parenthesized expression near the closing parenthesis. This still lets us expand tokens as we go, and avoids performance problems as long as there are enough parentheses. However, it would be better to avoid requiring the closing parenthesis to be present as soon as the corresponding opening parenthesis is read: the closing parenthesis may still be hidden in a macro yet to be expanded.

Hence, we need to go for some fine expansion control: the result is stored *before* the start!

Let us illustrate this idea in a simple model: adding positive integers which may be resulting from the expansion of macros, or may be values of registers. Assume that one number, say, 12345, has already been found, and that we want to parse the next number. The current status of the code may look as follows.

```

\exp_after:wN \add:ww \int_value:w 12345 \exp_after:wN ;
\exp:w \operand:w <stuff>

```

One step of expansion expands `\exp_after:wN`, which triggers the primitive `\int_value:w`, which reads the five digits we have already found, 12345. This integer is unfinished, causing the second `\exp_after:wN` to expand, and to trigger the construction

`\exp:w`, which expands `\operand:w`, defined to read what follows and make a number out of it, then leave `\exp_end:`, the number, and a semicolon in the input stream. Once `\operand:w` is done expanding, we obtain essentially

```
\exp_after:wN \add:ww \int_value:w 12345 ;
\exp:w \exp_end: 333444 ;
```

where in fact `\exp_after:wN` has already been expanded, `\int_value:w` has already seen 12345, and `\exp:w` is still looking for a number. It finds `\exp_end:`, hence expands to nothing. Now, `\int_value:w` sees the `;`, which cannot be part of a number. The expansion stops, and we are left with

```
\add:ww 12345 ; 333444 ;
```

which can safely perform the addition by grabbing two arguments delimited by `;`.

If we were to continue parsing the expression, then the following number should also be cleaned up before the next use of a binary operation such as `\add:ww`. Just like `\int_value:w 12345 \exp_after:wN ;` expanded what follows once, we need `\add:ww` to do the calculation, and in the process to expand the following once. This is also true in our real application: all the functions of the form `__fp..._o:ww` expand what follows once. This comes at the cost of leaving tokens in the input stack, and we need to be careful not to waste this memory. All of our discussion above is nice but simplistic, as operations should not simply be performed in the order they appear.

28.1.2 Precedence and infix operators

The various operators we will encounter have different precedences, which influence the order of calculations: $1 + 2 \times 3 = 1 + (2 \times 3)$ because \times has a higher precedence than $+$. The true analog of our macro `\operand:w` must thus take care of that. When looking for an operand, it needs to perform calculations until reaching an operator which has lower precedence than the one which called `\operand:w`. This means that `\operand:w` must know what the previous binary operator is, or rather, its precedence: we thus rename it `\operand:Nw`. Let us describe as an example how we plan to do the calculation $41 - 2^3 * 4 + 5$. More precisely we describe how to perform the first operation in this expression. Here, we abuse notations: the first argument of `\operand:Nw` should be an integer constant (`\c__fp_prec_plus_int, ...`) equal to the precedence of the given operator, not directly the operator itself.

- Clean up 41 and find $-$. We call `\operand:Nw -` to find the second operand.
- Clean up 2 and find \wedge .
- Compare the precedences of $-$ and \wedge . Since the latter is higher, we need to compute the exponentiation. For this, find the second operand with a nested call to `\operand:Nw \wedge`.
- Clean up 3 and find $*$.
- Compare the precedences of \wedge and $*$. Since the former is higher, `\operand:Nw \wedge` has found the second operand of the exponentiation, which is computed: $2^3 = 8$.
- We now have $41 - 8 * 4 + 5$, and `\operand:Nw -` is still looking for a second operand for the subtraction. Is it 8?

- Compare the precedences of $-$ and $*$. Since the latter is higher, we are not done with 8. Call `\operand:Nw *` to find the second operand of the multiplication.
- Clean up 4, and find $+$.
- Compare the precedences of $*$ and $+$. Since the former is higher, `\operand:Nw *` has found the second operand of the multiplication, which is computed: $8 * 4 = 32$.
- We now have $41 - 32 + 5$, and `\operand:Nw -` is still looking for a second operand for the subtraction. Is it 32?
- Compare the precedences of $-$ and $+$. Since they are equal, `\operand:Nw -` has found the second operand for the subtraction, which is computed: $41 - 32 = 9$.
- We now have $9 + 5$.

The procedure above stops short of performing all computations, but adding a surrounding call to `\operand:Nw` with a very low precedence ensures that all computations are performed before `\operand:Nw` is done. Adding a trailing marker with the same very low precedence prevents the surrounding `\operand:Nw` from going beyond the marker.

The pattern above to find an operand for a given operator, is to find one number and the next operator, then compare precedences to know if the next computation should be done. If it should, then perform it after finding its second operand, and look at the next operator, then compare precedences to know if the next computation should be done. This continues until we find that the next computation should not be done. Then, we stop.

We are now ready to get a bit more technical and describe which of the `l3fp-parse` functions correspond to each step above.

First, `__fp_parse_operand:Nw` is the `\operand:Nw` function above, with small modifications due to expansion issues discussed later. We denote by $\langle precedence \rangle$ the argument of `__fp_parse_operand:Nw`, that is, the precedence of the binary operator whose operand we are trying to find. The basic action is to read numbers from the input stream. This is done by `__fp_parse_one:Nw`. A first approximation of this function is that it reads one $\langle number \rangle$, performing no computation, and finds the following binary $\langle operator \rangle$. Then it expands to

$$\langle number \rangle \\ __fp_parse_infix_ \langle operator \rangle : N \langle precedence \rangle$$

expanding the `infix` auxiliary before leaving the above in the input stream.

We now explain the `infix` auxiliaries. We need some flexibility in how we treat the case of equal precedences: most often, the first operation encountered should be performed, such as $1 - 2 - 3$ being computed as $(1 - 2) - 3$, but 2^3^4 should be evaluated as $2^{(3^4)}$ instead. For this reason, and to support the equivalence between $**$ and \wedge more easily, each binary operator is converted to a control sequence `__fp_parse_infix_ \langle operator \rangle : N` when it is encountered for the first time. Instead of passing both precedences to a test function to do the comparison steps above, we pass the $\langle precedence \rangle$ (of the earlier operator) to the `infix` auxiliary for the following $\langle operator \rangle$, to know whether to perform the computation of the $\langle operator \rangle$. If it should not be performed, the `infix` auxiliary expands to

$$@ __use_none : n __fp_parse_infix_ \langle operator \rangle : N$$

and otherwise it calls `__fp_parse_operand:Nw` with the precedence of the $\langle operator \rangle$ to find its second operand $\langle number_2 \rangle$ and the next $\langle operator_2 \rangle$, and expands to

```
@ \__fp_parse_apply_binary:NwNwN
  \langle operator \rangle \langle number_2 \rangle
@ \__fp_parse_infix_\langle operator_2 \rangle:N
```

The `infix` function is responsible for comparing precedences, but cannot directly call the computation functions, because the first operand $\langle number \rangle$ is before the `infix` function in the input stream. This is why we stop the expansion here and give control to another function to close the loop.

A definition of `__fp_parse_operand:Nw` $\langle precedence \rangle$ with some of the expansion control removed is

```
\exp_after:wN \__fp_parse_continue:NwN
\exp_after:wN \langle precedence \rangle
\exp:w \exp_end_continue_f:w
  \__fp_parse_one:Nw \langle precedence \rangle
```

This expands `__fp_parse_one:Nw` $\langle precedence \rangle$ completely, which finds a number, wraps the next $\langle operator \rangle$ into an `infix` function, feeds this function the $\langle precedence \rangle$, and expands it, yielding either

```
\__fp_parse_continue:NwN \langle precedence \rangle
\langle number \rangle @
\use_none:n \__fp_parse_infix_\langle operator \rangle:N
```

or

```
\__fp_parse_continue:NwN \langle precedence \rangle
\langle number \rangle @
\__fp_parse_apply_binary:NwNwN
  \langle operator \rangle \langle number_2 \rangle
@ \__fp_parse_infix_\langle operator_2 \rangle:N
```

The definition of `__fp_parse_continue:NwN` is then very simple:

```
\cs_new:Npn \__fp_parse_continue:NwN #1#2@#3 { #3 #1 #2 @ }
```

In the first case, `#3` is `\use_none:n`, yielding

```
\use_none:n \langle precedence \rangle \langle number \rangle @
\__fp_parse_infix_\langle operator \rangle:N
```

then $\langle number \rangle @ __fp_parse_infix_ \langle operator \rangle:N$. In the second case, `#3` is `__fp_parse_apply_binary:NwNwN`, whose role is to compute $\langle number \rangle \langle operator \rangle \langle number_2 \rangle$ and to prepare for the next comparison of precedences: first we get

```
\__fp_parse_apply_binary:NwNwN
  \langle precedence \rangle \langle number \rangle @
  \langle operator \rangle \langle number_2 \rangle
@ \__fp_parse_infix_\langle operator_2 \rangle:N
```

then

```

\exp_after:wN \__fp_parse_continue:NwN
\exp_after:wN <precedence>
\exp:w \exp_end_continue_f:w
\__fp_<operator>_o:ww <number> <number2>
\exp:w \exp_end_continue_f:w
\__fp_parse_infix_<operator2>:N <precedence>

```

where `__fp_<operator>_o:ww` computes `<number> <operator> <number2>` and expands after the result, thus triggers the comparison of the precedence of the `<operator2>` and the `<precedence>`, continuing the loop.

We have introduced the most important functions here, and the next few paragraphs we describe various subtleties.

28.1.3 Prefix operators, parentheses, and functions

Prefix operators (unary `-`, `+`, `!`) and parentheses are taken care of by the same mechanism, and functions (`sin`, `exp`, etc.) as well. Finding the argument of the unary `-`, for instance, is very similar to grabbing the second operand of a binary infix operator, with a subtle precedence explained below. Once that operand is found, the operator can be applied to it (for the unary `-`, this simply flips the sign). A left parenthesis is just a prefix operator with a very low precedence equal to that of the closing parenthesis (which is treated as an infix operator, since it normally appears just after numbers), so that all computations are performed until the closing parenthesis. The prefix operator associated to the left parenthesis does not alter its argument, but it removes the closing parenthesis (with some checks).

Prefix operators are the reason why we only summarily described the function `__fp_parse_one:Nw` earlier. This function is responsible for reading in the input stream the first possible `<number>` and the next infix `<operator>`. If what follows `__fp_parse_one:Nw <precedence>` is a prefix operator, then we must find the operand of this prefix operator through a nested call to `__fp_parse_operand:Nw` with the appropriate precedence, then apply the operator to the operand found to yield the result of `__fp_parse_one:Nw`. So far, all is simple.

The unary operators `+`, `-`, `!` complicate things a little bit: `-3**2` should be $-(3^2) = -9$, and not $(-3)^2 = 9$. This would easily be done by giving `-` a lower precedence, equal to that of the infix `+` and `-`. Unfortunately, this fails in cases such as `3**-2*4`, yielding $3^{-2 \times 4}$ instead of the correct $3^{-2} \times 4$. A second attempt would be to call `__fp_parse_operand:Nw` with the `<precedence>` of the previous operator, but `0>-2+3` is then parsed as `0>-(2+3)`: the addition is performed because it binds more tightly than the comparison which precedes `-`. The correct approach is for a unary `-` to perform operations whose precedence is greater than both that of the previous operation, and that of the unary `-` itself. The unary `-` is given a precedence higher than multiplication and division. This does not lead to any surprising result, since $-(x/y) = (-x)/y$ and similarly for multiplication, and it reduces the number of nested calls to `__fp_parse_operand:Nw`.

Functions are implemented as prefix operators with very high precedence, so that their argument is the first number that can possibly be built.

Note that contrarily to the `infix` functions discussed earlier, the `prefix` functions do perform tests on the previous `<precedence>` to decide whether to find an argument or not, since we know that we need a number, and must never stop there.

28.1.4 Numbers and reading tokens one by one

So far, we have glossed over one important point: what is a “number”? A number is typically given in the form $\langle\textit{significand}\rangle\mathbf{e}\langle\textit{exponent}\rangle$, where the $\langle\textit{significand}\rangle$ is any non-empty string composed of decimal digits and at most one decimal separator (a period), the exponent “ $\mathbf{e}\langle\textit{exponent}\rangle$ ” is optional and is composed of an exponent mark **e** followed by a possibly empty string of signs + or - and a non-empty string of decimal digits. The $\langle\textit{significand}\rangle$ can also be an integer, dimension, skip, or muskip variable, in which case dimensions are converted from points (or mu units) to floating points, and the $\langle\textit{exponent}\rangle$ can also be an integer variable. Numbers can also be given as floating point variables, or as named constants such as **nan**, **inf** or **pi**. We may add more types in the future.

When `__fp_parse_one:Nw` is looking for a “number”, here is what happens.

- If the next token is a control sequence with the meaning of `\scan_stop:`, it can be: `\s__fp`, in which case our job is done, as what follows is an internal floating point number, or `\s__fp_expr_mark`, in which case the expression has come to an early end, as we are still looking for a number here, or something else, in which case we consider the control sequence to be a bad variable resulting from `c`-expansion.
- If the next token is a control sequence with a different meaning, we assume that it is a register, unpack it with `\tex_the:D`, and use its value (in `pt` for dimensions and skips, `mu` for muskips) as the $\langle\textit{significand}\rangle$ of a number: we look for an exponent.
- If the next token is a digit, we remove any leading zeros, then read a significand larger than 1 if the next character is a digit, read a significand smaller than 1 if the next character is a period, or we have found a significand equal to 0 otherwise, and look for an exponent.
- If the next token is a letter, we collect more letters until the first non-letter: the resulting word may denote a function such as `asin`, a constant such as `pi` or be unknown. In the first case, we call `__fp_parse_operand:Nw` to find the argument of the function, then apply the function, before declaring that we are done. Otherwise, we are done, either with the value of the constant, or with the value **nan** for unknown words.
- If the next token is anything else, we check whether it is a known prefix operator, in which case `__fp_parse_operand:Nw` finds its operand. If it is not known, then either a number is missing (if the token is a known infix operator) or the token is simply invalid in floating point expressions.

Once a number is found, `__fp_parse_one:Nw` also finds an infix operator. This goes as follows.

- If the next token is a control sequence, it could be the special marker `\s__fp_expr_mark`, and otherwise it is a case of juxtaposing numbers, such as `2\c_zero_int`, with an implied multiplication.
- If the next token is a letter, it is also a case of juxtaposition, as letters cannot be proper infix operators.
- Otherwise (including in the case of digits), if the token is a known infix operator, the appropriate `__fp_infix_<operator>:N` function is built, and if it does not exist, we complain. In particular, the juxtaposition `\c_zero_int 2` is disallowed.

In the above, we need to test whether a character token #1 is a digit:

```
\if_int_compare:w 9 < 1 \token_to_str:N #1 \exp_stop_f:
  is a digit
\else:
  not a digit
\fi:
```

To exclude 0, replace 9 by 10. The use of `\token_to_str:N` ensures that a digit with any catcode is detected. To test if a character token is a letter, we need to work with its character code, testing if ‘#1 lies in [65,90] (uppercase letters) or [97,112] (lowercase letters)

```
\if_int_compare:w \__fp_int_eval:w
  ( ‘#1 \if_int_compare:w ‘#1 > ‘Z - 32 \fi: ) / 26 = 3 \exp_stop_f:
  is a letter
\else:
  not a letter
\fi:
```

At all steps, we try to accept all category codes: when #1 is kept to be used later, it is almost always converted to category code other through `\token_to_str:N`. More precisely, catcodes {3,6,7,8,11,12} should work without trouble, but not {1,2,4,10,13}, and of course {0,5,9} cannot become tokens.

Floating point expressions should behave as much as possible like ε -TeX-based integer expressions and dimension expressions. In particular, `f`-expansion should be performed as the expression is read, token by token, forcing the expansion of protected macros, and ignoring spaces. One advantage of expanding at every step is that restricted expandable functions can then be used in floating point expressions just as they can be in other kinds of expressions. Problematically, spaces stop `f`-expansion: for instance, the macro `\X` below would not be expanded if we simply performed `f`-expansion.

```
\DeclareDocumentCommand {\test} {m} { \fp_eval:n {#1} }
\ExplSyntaxOff
\test { 1 + \X }
```

Of course, spaces typically do not appear in a code setting, but may very easily come in document-level input, from which some expressions may come. To avoid this problem, at every step, we do essentially what `\use:f` would do: take an argument, put it back in the input stream, then `f`-expand it. This is not a complete solution, since a macro’s expansion could contain leading spaces which would stop the `f`-expansion before further macro calls are performed. However, in practice it should be enough: in particular, floating point numbers are correctly expanded to the underlying `\s__fp ...` structure. The `f`-expansion is performed by `__fp_parse_expand:w`.

28.2 Main auxiliary functions

```
\__fp_parse_operand:Nw \exp:w \__fp_parse_operand:Nw <precedence> \__fp_parse_expand:w
Reads the "...", performing every computation with a precedence higher than
<precedence>, then expands to

<result> @ \__fp_parse_infix_<operation>:N ...
```

where the $\langle operation \rangle$ is the first operation with a lower precedence, possibly `end`, and the “...” start just after the $\langle operation \rangle$.

(End definition for `_fp_parse_operand:Nw`.)

```
\_fp_parse_infix_+:N      \_fp_parse_infix_+:N  $\langle precedence \rangle$  ...
                          If + has a precedence higher than the  $\langle precedence \rangle$ , cleans up a second  $\langle operand \rangle$  and
                          finds the  $\langle operation_2 \rangle$  which follows, and expands to

                          @ \_fp_parse_apply_binary:NwNwN +  $\langle operand \rangle$  @ \_fp_parse_infix_ $\langle operation_2 \rangle$ :N
                          ...
```

Otherwise expands to

```
@ \use_none:n \_fp_parse_infix_+:N ...
```

A similar function exists for each infix operator.

(End definition for `_fp_parse_infix_+:N`.)

```
\_fp_parse_one:Nw      \_fp_parse_one:Nw  $\langle precedence \rangle$  ...
                        Cleans up one or two operands depending on how the precedence of the next oper-
                        ation compares to the  $\langle precedence \rangle$ . If the following  $\langle operation \rangle$  has a precedence higher
                        than  $\langle precedence \rangle$ , expands to
```

```
 $\langle operand_1 \rangle$  @ \_fp_parse_apply_binary:NwNwN  $\langle operation \rangle$   $\langle operand_2 \rangle$  @
\_fp_parse_infix_ $\langle operation_2 \rangle$ :N ...
```

and otherwise expands to

```
 $\langle operand \rangle$  @ \use_none:n \_fp_parse_infix_ $\langle operation \rangle$ :N ...
```

(End definition for `_fp_parse_one:Nw`.)

28.3 Helpers

```
\_fp_parse_expand:w      \exp:w \_fp_parse_expand:w  $\langle tokens \rangle$ 
```

This function must always come within a `\exp:w` expansion. The $\langle tokens \rangle$ should be the part of the expression that we have not yet read. This requires in particular closing all conditionals properly before expanding.

```
17150 \cs_new:Npn \_fp_parse_expand:w #1 { \exp_end_continue_f:w #1 }
```

(End definition for `_fp_parse_expand:w`.)

```
\_fp_parse_return_semicolon:w This very odd function swaps its position with the following \fi: and removes \_fp_
parse_expand:w normally responsible for expansion. That turns out to be useful.
```

```
17151 \cs_new:Npn \_fp_parse_return_semicolon:w
17152   #1 \fi: \_fp_parse_expand:w { \fi: ; #1 }
```

(End definition for `_fp_parse_return_semicolon:w`.)

```
\_fp_parse_digits_vii:N These functions must be called within an \int_value:w or \_fp_int_eval:w construc-
\_fp_parse_digits_vi:N tion. The first token which follows must be f-expanded prior to calling those functions.
\_fp_parse_digits_v:N The functions read tokens one by one, and output digits into the input stream, until
\_fp_parse_digits_iv:N meeting a non-digit, or up to a number of digits equal to their index. The full expansion
\_fp_parse_digits_iii:N is
```

```
\_fp_parse_digits_ii:N
\_fp_parse_digits_i:N
\_fp_parse_digits_:N
```

$\langle \text{digits} \rangle$; $\langle \text{filling } 0 \rangle$; $\langle \text{length} \rangle$

where $\langle \text{filling } 0 \rangle$ is a string of zeros such that $\langle \text{digits} \rangle \langle \text{filling } 0 \rangle$ has the length given by the index of the function, and $\langle \text{length} \rangle$ is the number of zeros in the $\langle \text{filling } 0 \rangle$ string. Each function puts a digit into the input stream and calls the next function, until we find a non-digit. We are careful to pass the tested tokens through $\backslash \text{token_to_str:N}$ to normalize their category code.

```

17153 \cs_set_protected:Npn \__fp_tmp:w #1 #2 #3
17154 {
17155   \cs_new:cpn { __fp_parse_digits_ #1 :N } ##1
17156   {
17157     \if_int_compare:w 9 < 1 \token_to_str:N ##1 \exp_stop_f:
17158     \token_to_str:N ##1 \exp_after:wN #2 \exp:w
17159     \else:
17160       \__fp_parse_return_semicolon:w #3 ##1
17161     \fi:
17162     \__fp_parse_expand:w
17163   }
17164 }
17165 \__fp_tmp:w {vii} \__fp_parse_digits_vi:N { 0000000 ; 7 }
17166 \__fp_tmp:w {vi} \__fp_parse_digits_v:N { 000000 ; 6 }
17167 \__fp_tmp:w {v} \__fp_parse_digits_iv:N { 00000 ; 5 }
17168 \__fp_tmp:w {iv} \__fp_parse_digits_iii:N { 0000 ; 4 }
17169 \__fp_tmp:w {iii} \__fp_parse_digits_ii:N { 000 ; 3 }
17170 \__fp_tmp:w {ii} \__fp_parse_digits_i:N { 00 ; 2 }
17171 \__fp_tmp:w {i} \__fp_parse_digits_:N { 0 ; 1 }
17172 \cs_new:Npn \__fp_parse_digits_:N { ; ; 0 }

```

(End definition for $\backslash \text{__fp_parse_digits_vii:N}$ and others.)

28.4 Parsing one number

$\backslash \text{__fp_parse_one:Nw}$ This function finds one number, and packs the symbol which follows in an $\backslash \text{__fp_parse_infix_...}$ csname. #1 is the previous $\langle \text{precedence} \rangle$, and #2 the first token of the operand. We distinguish four cases: #2 is equal to $\backslash \text{scan_stop:}$ in meaning, #2 is a different control sequence, #2 is a digit, and #2 is something else (this last case is split further later). Despite the earlier f-expansion, #2 may still be expandable if it was protected by $\backslash \text{exp_not:N}$, as may happen with the L^AT_EX 2_ε command $\backslash \text{protect}$. Using a well placed $\backslash \text{reverse_if:N}$, this case is sent to $\backslash \text{__fp_parse_one_fp:NN}$ which deals with it robustly.

```

17173 \cs_new:Npn \__fp_parse_one:Nw #1 #2
17174 {
17175   \if_catcode:w \scan_stop: \exp_not:N #2
17176   \exp_after:wN \if_meaning:w \exp_not:N #2 #2 \else:
17177     \exp_after:wN \reverse_if:N
17178   \fi:
17179   \if_meaning:w \scan_stop: #2
17180   \exp_after:wN \exp_after:wN
17181   \exp_after:wN \__fp_parse_one_fp:NN
17182   \else:
17183     \exp_after:wN \exp_after:wN
17184     \exp_after:wN \__fp_parse_one_register:NN
17185   \fi:

```

```

17186 \else:
17187 \if_int_compare:w 9 < 1 \token_to_str:N #2 \exp_stop_f:
17188 \exp_after:wN \exp_after:wN
17189 \exp_after:wN \__fp_parse_one_digit:NN
17190 \else:
17191 \exp_after:wN \exp_after:wN
17192 \exp_after:wN \__fp_parse_one_other:NN
17193 \fi:
17194 \fi:
17195 #1 #2
17196 }

```

(End definition for `__fp_parse_one:Nw`.)

```

\__fp_parse_one_fp:NN
\__fp_exp_after_expr_mark_f:nw
\__fp_exp_after_?_f:nw

```

This function receives a $\langle precedence \rangle$ and a control sequence equal to `\scan_stop:` in meaning. There are three cases.

- `\s__fp` starts a floating point number, and we call `__fp_exp_after_f:nw`, which f-expands after the floating point.
- `\s__fp_expr_mark` is a premature end, we call `__fp_exp_after_expr_mark_f:nw`, which triggers an `fp-early-end` error.
- For a control sequence not containing `\s__fp`, we call `__fp_exp_after_?_f:nw`, causing a `bad-variable` error.

This scheme is extensible: additional types can be added by starting the variables with a scan mark of the form `\s__fp_⟨type⟩` and defining `__fp_exp_after_⟨type⟩_f:nw`. In all cases, we make sure that the second argument of `__fp_parse_infix:NN` is correctly expanded. A special case only enabled in L^AT_EX 2_ε is that if `\protect` is encountered then the error message mentions the control sequence which follows it rather than `\protect` itself. The test for L^AT_EX 2_ε uses `\@unexpandable@protect` rather than `\protect` because `\protect` is often `\scan_stop:` hence “does not exist”.

```

17197 \cs_new:Npn \__fp_parse_one_fp:NN #1
17198 {
17199 \__fp_exp_after_any_f:nw
17200 {
17201 \exp_after:wN \__fp_parse_infix:NN
17202 \exp_after:wN #1 \exp:w \__fp_parse_expand:w
17203 }
17204 }
17205 \cs_new:Npn \__fp_exp_after_expr_mark_f:nw #1
17206 {
17207 \int_case:nnF { \exp_after:wN \use_i:nnn \use_none:nnn #1 }
17208 {
17209 \c__fp_prec_comma_int { }
17210 \c__fp_prec_tuple_int { }
17211 \c__fp_prec_end_int
17212 {
17213 \exp_after:wN \c__fp_empty_tuple_fp
17214 \exp:w \exp_end_continue_f:w
17215 }
17216 }
17217 {

```



```

17218     \_kernel_msg_expandable_error:nn { kernel } { fp-early-end }
17219     \exp_after:wN \c_nan_fp \exp:w \exp_end_continue_f:w
17220   }
17221   #1
17222 }
17223 \cs_new:cpn { __fp_exp_after_?_f:nw } #1#2
17224 {
17225   \_kernel_msg_expandable_error:nnn { kernel } { bad-variable }
17226   {#2}
17227   \exp_after:wN \c_nan_fp \exp:w \exp_end_continue_f:w #1
17228 }
17229 \cs_set_protected:Npn \__fp_tmp:w #1
17230 {
17231   \cs_if_exist:NT #1
17232   {
17233     \cs_gset:cpn { __fp_exp_after_?_f:nw } ##1##2
17234     {
17235       \exp_after:wN \c_nan_fp \exp:w \exp_end_continue_f:w ##1
17236       \str_if_eq:nnTF {##2} { \protect }
17237       {
17238         \cs_if_eq:NNTF ##2 #1 { \use_i:nn } { \use:n }
17239         {
17240           \_kernel_msg_expandable_error:nnn { kernel }
17241           { fp-robust-cmd }
17242         }
17243       }
17244       {
17245         \_kernel_msg_expandable_error:nnn { kernel }
17246         { bad-variable } {##2}
17247       }
17248     }
17249   }
17250 }
17251 \exp_args:Nc \__fp_tmp:w { @unexpandable@protect }

```

(End definition for `__fp_parse_one_fp:NN`, `__fp_exp_after_expr_mark_f:nw`, and `__fp_exp_after_?_f:nw`.)

```

\__fp_parse_one_register:NN
  \__fp_parse_one_register_aux:Nw
  \__fp_parse_one_register_auxii:wwwNw
  \__fp_parse_one_register_int:www
  \__fp_parse_one_register_mu:www
  \__fp_parse_one_register_dim:ww

```

This is called whenever #2 is a control sequence other than `\scan_stop:` in meaning. We special-case `\wd`, `\ht`, `\dp` (see later) and otherwise assume that it is a register, but carefully unpack it with `\tex_the:D` within braces. First, we find the exponent following #2. Then we unpack #2 with `\tex_the:D`, and the `auxii` auxiliary distinguishes integer registers from dimensions/skips from muskips, according to the presence of a period and/or of `pt`. For integers, simply convert $\langle value \rangle e \langle exponent \rangle$ to a floating point number with `__fp_parse:n` (this is somewhat wasteful). For other registers, the decimal rounding provided by $\mathrm{T}_{\mathrm{E}}\mathrm{X}$ does not accurately represent the binary value that it manipulates, so we extract this binary value as a number of scaled points with `\int_value:w \dim_to_decimal_in_sp:n { $\langle decimal value \rangle$ pt }`, and use an auxiliary of `\dim_to_fp:n`, which performs the multiplication by 2^{-16} , correctly rounded.

```

17252 \cs_new:Npn \__fp_parse_one_register:NN #1#2
17253 {
17254   \exp_after:wN \__fp_parse_infix_after_operand:NwN
17255   \exp_after:wN #1
17256   \exp:w \exp_end_continue_f:w

```

```

17257 \__fp_parse_one_register_special:N #2
17258 \exp_after:wN \__fp_parse_one_register_aux:Nw
17259 \exp_after:wN #2
17260 \int_value:w
17261 \exp_after:wN \__fp_parse_exponent:N
17262 \exp:w \__fp_parse_expand:w
17263 }
17264 \cs_new:Npx \__fp_parse_one_register_aux:Nw #1
17265 {
17266 \exp_not:n
17267 {
17268 \exp_after:wN \use:nn
17269 \exp_after:wN \__fp_parse_one_register_auxii:wwwNw
17270 }
17271 \exp_not:N \exp_after:wN { \exp_not:N \tex_the:D #1 }
17272 ; \exp_not:N \__fp_parse_one_register_dim:ww
17273 \tl_to_str:n { pt } ; \exp_not:N \__fp_parse_one_register_mu:www
17274 . \tl_to_str:n { pt } ; \exp_not:N \__fp_parse_one_register_int:www
17275 \s__fp_stop
17276 }
17277 \exp_args:Nno \use:nn
17278 { \cs_new:Npn \__fp_parse_one_register_auxii:wwwNw #1 . #2 }
17279 { \tl_to_str:n { pt } #3 ; #4#5 \s__fp_stop }
17280 { #4 #1.#2; }
17281 \exp_args:Nno \use:nn
17282 { \cs_new:Npn \__fp_parse_one_register_mu:www #1 }
17283 { \tl_to_str:n { mu } ; #2 ; }
17284 { \__fp_parse_one_register_dim:ww #1 ; }
17285 \cs_new:Npn \__fp_parse_one_register_int:www #1; #2.; #3;
17286 { \__fp_parse:n { #1 e #3 } }
17287 \cs_new:Npn \__fp_parse_one_register_dim:ww #1; #2;
17288 {
17289 \exp_after:wN \__fp_from_dim_test:ww
17290 \int_value:w #2 \exp_after:wN ,
17291 \int_value:w \dim_to_decimal_in_sp:n { #1 pt } ;
17292 }

```

(End definition for `__fp_parse_one_register:NN` and others.)

```

\__fp_parse_one_register_special:N
\__fp_parse_one_register_math:NNw
\__fp_parse_one_register_wd:w
\__fp_parse_one_register_wd:Nw

```

The `\wd`, `\dp`, `\ht` primitives expect an integer argument. We abuse the exponent parser to find the integer argument: simply include the exponent marker `e`. Once that “exponent” is found, use `\tex_the:D` to find the box dimension and then copy what we did for dimensions.

```

17293 \cs_new:Npn \__fp_parse_one_register_special:N #1
17294 {
17295 \if_meaning:w \box_wd:N #1 \__fp_parse_one_register_wd:w \fi:
17296 \if_meaning:w \box_ht:N #1 \__fp_parse_one_register_wd:w \fi:
17297 \if_meaning:w \box_dp:N #1 \__fp_parse_one_register_wd:w \fi:
17298 \if_meaning:w \infty #1
17299 \__fp_parse_one_register_math:NNw \infty #1
17300 \fi:
17301 \if_meaning:w \pi #1
17302 \__fp_parse_one_register_math:NNw \pi #1
17303 \fi:

```

```

17304 }
17305 \cs_new:Npn \__fp_parse_one_register_math:NNw
17306   #1#2#3#4 \__fp_parse_expand:w
17307   {
17308     #3
17309     \str_if_eq:nnTF {#1} {#2}
17310     {
17311       \__kernel_msg_expandable_error:nnn
17312       { kernel } { fp-infty-pi } {#1}
17313       \c_nan_fp
17314     }
17315     { #4 \__fp_parse_expand:w }
17316   }
17317 \cs_new:Npn \__fp_parse_one_register_wd:w
17318   #1#2 \exp_after:wN #3#4 \__fp_parse_expand:w
17319   {
17320     #1
17321     \exp_after:wN \__fp_parse_one_register_wd:Nw
17322     #4 \__fp_parse_expand:w e
17323   }
17324 \cs_new:Npn \__fp_parse_one_register_wd:Nw #1#2 ;
17325   {
17326     \exp_after:wN \__fp_from_dim_test:ww
17327     \exp_after:wN 0 \exp_after:wN ,
17328     \int_value:w \dim_to_decimal_in_sp:n { #1 #2 } ;
17329   }

```

(End definition for __fp_parse_one_register_special:N and others.)

__fp_parse_one_digit:NN A digit marks the beginning of an explicit floating point number. Once the number is found, we catch the case of overflow and underflow with __fp_sanitize:wN, then __fp_parse_infix_after_operand:NwN expands __fp_parse_infix:NN after the number we find, to wrap the following infix operator as required. Finding the number itself begins by removing leading zeros: further steps are described later.

```

17330 \cs_new:Npn \__fp_parse_one_digit:NN #1
17331   {
17332     \exp_after:wN \__fp_parse_infix_after_operand:NwN
17333     \exp_after:wN #1
17334     \exp:w \exp_end_continue_f:w
17335     \exp_after:wN \__fp_sanitize:wN
17336     \int_value:w \__fp_int_eval:w 0 \__fp_parse_trim_zeros:N
17337   }

```

(End definition for __fp_parse_one_digit:NN.)

__fp_parse_one_other:NN For this function, #2 is a character token which is not a digit. If it is an ASCII letter, __fp_parse_letters:N beyond this one and give the result to __fp_parse_word:Nw. Otherwise, the character is assumed to be a prefix operator, and we build __fp_parse_prefix_{operator}:Nw.

```

17338 \cs_new:Npn \__fp_parse_one_other:NN #1 #2
17339   {
17340     \if_int_compare:w
17341       \__fp_int_eval:w
17342       ( '#2 \if_int_compare:w '#2 > 'Z - 32 \fi: ) / 26

```

```

17343         = 3 \exp_stop_f:
17344         \exp_after:wN \__fp_parse_word:Nw
17345         \exp_after:wN #1
17346         \exp_after:wN #2
17347         \exp:w \exp_after:wN \__fp_parse_letters:N
17348         \exp:w
17349     \else:
17350         \exp_after:wN \__fp_parse_prefix:NNN
17351         \exp_after:wN #1
17352         \exp_after:wN #2
17353         \cs:w
17354         __fp_parse_prefix_ \token_to_str:N #2 :Nw
17355         \exp_after:wN
17356         \cs_end:
17357         \exp:w
17358     \fi:
17359     \__fp_parse_expand:w
17360 }

```

(End definition for __fp_parse_one_other:NN.)

__fp_parse_word:Nw
__fp_parse_letters:N

Finding letters is a simple recursion. Once __fp_parse_letters:N has done its job, we try to build a control sequence from the word #2. If it is a known word, then the corresponding action is taken, and otherwise, we complain about an unknown word, yield \c_nan_fp, and look for the following infix operator. Note that the unknown word could be a mistyped function as well as a mistyped constant, so there is no way to tell whether to look for arguments; we do not. The standard requires “inf” and “infinity” and “nan” to be recognized regardless of case, but we probably don’t want to allow every l3fp word to have an arbitrary mixture of lower and upper case, so we test and use a differently-named control sequence.

```

17361 \cs_new:Npn \__fp_parse_word:Nw #1#2;
17362 {
17363     \cs_if_exist_use:cF { __fp_parse_word_#2:N }
17364     {
17365         \cs_if_exist_use:cF
17366         { __fp_parse_caseless_ \str_foldcase:n {#2} :N }
17367         {
17368             \__kernel_msg_expandable_error:nnn
17369             { kernel } { unknown-fp-word } {#2}
17370             \exp_after:wN \c_nan_fp \exp:w \exp_end_continue_f:w
17371             \__fp_parse_infix:NN
17372         }
17373     }
17374     #1
17375 }
17376 \cs_new:Npn \__fp_parse_letters:N #1
17377 {
17378     \exp_end_continue_f:w
17379     \if_int_compare:w
17380     \if_catcode:w \scan_stop: \exp_not:N #1
17381     0
17382     \else:
17383     \__fp_int_eval:w
17384     ( '#1 \if_int_compare:w '#1 > 'Z - 32 \fi: ) / 26

```

```

17385     \fi:
17386     = 3 \exp_stop_f:
17387     \exp_after:wN #1
17388     \exp:w \exp_after:wN \_fp_parse_letters:N
17389     \exp:w
17390   \else:
17391     \_fp_parse_return_semicolon:w #1
17392   \fi:
17393   \_fp_parse_expand:w
17394 }

```

(End definition for _fp_parse_word:Nw and _fp_parse_letters:N.)

```

\_fp_parse_prefix:NNN
\_fp_parse_prefix_unknown:NNN

```

For this function, #1 is the previous $\langle precedence \rangle$, #2 is the operator just seen, and #3 is a control sequence which implements the operator if it is a known operator. If this control sequence is `\scan_stop:`, then the operator is in fact unknown. Either the expression is missing a number there (if the operator is valid as an infix operator), and we put `nan`, wrapping the infix operator in a csname as appropriate, or the character is simply invalid in floating point expressions, and we continue looking for a number, starting again from `_fp_parse_one:Nw`.

```

17395 \cs_new:Npn \_fp_parse_prefix:NNN #1#2#3
17396 {
17397   \if_meaning:w \scan_stop: #3
17398   \exp_after:wN \_fp_parse_prefix_unknown:NNN
17399   \exp_after:wN #2
17400   \fi:
17401   #3 #1
17402 }
17403 \cs_new:Npn \_fp_parse_prefix_unknown:NNN #1#2#3
17404 {
17405   \cs_if_exist:cTF { \_fp_parse_infix_ \token_to_str:N #1 :N }
17406   {
17407     \_kernel_msg_expandable_error:nnn
17408     { kernel } { fp-missing-number } {#1}
17409     \exp_after:wN \c_nan_fp \exp:w \exp_end_continue_f:w
17410     \_fp_parse_infix:NN #3 #1
17411   }
17412   {
17413     \_kernel_msg_expandable_error:nnn
17414     { kernel } { fp-unknown-symbol } {#1}
17415     \_fp_parse_one:Nw #3
17416   }
17417 }

```

(End definition for _fp_parse_prefix:NNN and _fp_parse_prefix_unknown:NNN.)

28.4.1 Numbers: trimming leading zeros

Numbers are parsed as follows: first we trim leading zeros, then if the next character is a digit, start reading a significand ≥ 1 with the set of functions `_fp_parse_large...`; if it is a period, the significand is < 1 ; and otherwise it is zero. In the second case, trim additional zeros after the period, counting them for an exponent shift $\langle exp_1 \rangle < 0$, then read the significand with the set of functions `_fp_parse_small...`. Once the significand is read, read the exponent if `e` is present.

`__fp_parse_trim_zeros:N` This function expects an already expanded token. It removes any leading zero, then distinguishes three cases: if the first non-zero token is a digit, then call `__fp_parse_large:N` (the significand is ≥ 1); if it is `.`, then continue trimming zeros with `__fp_parse_strim_zeros:N`; otherwise, our number is exactly zero, and we call `__fp_parse_zero:` to take care of that case.

```

17418 \cs_new:Npn \__fp_parse_trim_zeros:N #1
17419 {
17420   \if:w 0 \exp_not:N #1
17421     \exp_after:wN \__fp_parse_trim_zeros:N
17422     \exp:w
17423   \else:
17424     \if:w . \exp_not:N #1
17425       \exp_after:wN \__fp_parse_strim_zeros:N
17426       \exp:w
17427     \else:
17428       \__fp_parse_trim_end:w #1
17429     \fi:
17430   \fi:
17431   \__fp_parse_expand:w
17432 }
17433 \cs_new:Npn \__fp_parse_trim_end:w #1 \fi: \fi: \__fp_parse_expand:w
17434 {
17435   \fi:
17436   \fi:
17437   \if_int_compare:w 9 < 1 \token_to_str:N #1 \exp_stop_f:
17438     \exp_after:wN \__fp_parse_large:N
17439   \else:
17440     \exp_after:wN \__fp_parse_zero:
17441   \fi:
17442   #1
17443 }
```

(End definition for `__fp_parse_trim_zeros:N` and `__fp_parse_trim_end:w`.)

`__fp_parse_strim_zeros:N` If we have removed all digits until a period (or if the body started with a period), then enter the “small_trim” loop which outputs `-1` for each removed 0. Those `-1` are added to an integer expression waiting for the exponent. If the first non-zero token is a digit, call `__fp_parse_small:N` (our significand is smaller than 1), and otherwise, the number is an exact zero. The name `strim` stands for “small trim”.

```

17444 \cs_new:Npn \__fp_parse_strim_zeros:N #1
17445 {
17446   \if:w 0 \exp_not:N #1
17447     - 1
17448     \exp_after:wN \__fp_parse_strim_zeros:N \exp:w
17449   \else:
17450     \__fp_parse_strim_end:w #1
17451   \fi:
17452   \__fp_parse_expand:w
17453 }
17454 \cs_new:Npn \__fp_parse_strim_end:w #1 \fi: \__fp_parse_expand:w
17455 {
17456   \fi:
17457   \if_int_compare:w 9 < 1 \token_to_str:N #1 \exp_stop_f:
```

```

17458     \exp_after:wN \__fp_parse_small:N
17459   \else:
17460     \exp_after:wN \__fp_parse_zero:
17461   \fi:
17462   #1
17463 }

```

(End definition for __fp_parse_strim_zeros:N and __fp_parse_strim_end:w.)

__fp_parse_zero: After reading a significand of 0, find any exponent, then put a sign of 1 for __fp-sanitize:wN, which removes everything and leaves an exact zero.

```

17464 \cs_new:Npn \__fp_parse_zero:
17465 {
17466   \exp_after:wN ; \exp_after:wN 1
17467   \int_value:w \__fp_parse_exponent:N
17468 }

```

(End definition for __fp_parse_zero:.)

28.4.2 Number: small significand

__fp_parse_small:N This function is called after we have passed the decimal separator and removed all leading zeros from the significand. It is followed by a non-zero digit (with any catcode). The goal is to read up to 16 digits. But we can't do that all at once, because \int_value:w (which allows us to collect digits and continue expanding) can only go up to 9 digits. Hence we grab digits in two steps of 8 digits. Since #1 is a digit, read seven more digits using __fp_parse_digits_vii:N. The small_leading auxiliary leaves those digits in the \int_value:w, and grabs some more, or stops if there are no more digits. Then the pack_leading auxiliary puts the various parts in the appropriate order for the processing further up.

```

17469 \cs_new:Npn \__fp_parse_small:N #1
17470 {
17471   \exp_after:wN \__fp_parse_pack_leading:NNNNNww
17472   \int_value:w \__fp_int_eval:w 1 \token_to_str:N #1
17473   \exp_after:wN \__fp_parse_small_leading:wwNN
17474   \int_value:w 1
17475   \exp_after:wN \__fp_parse_digits_vii:N
17476   \exp:w \__fp_parse_expand:w
17477 }

```

(End definition for __fp_parse_small:N.)

__fp_parse_small_leading:wwNN __fp_parse_small_leading:wwNN 1 <digits> ; <zeros> ; <number of zeros>

We leave <digits> <zeros> in the input stream: the functions used to grab digits are such that this constitutes digits 1 through 8 of the significand. Then prepare to pack 8 more digits, with an exponent shift of zero (this shift is used in the case of a large significand). If #4 is a digit, leave it behind for the packing function, and read 6 more digits to reach a total of 15 digits: further digits are involved in the rounding. Otherwise put 8 zeros in to complete the significand, then look for an exponent.

```

17478 \cs_new:Npn \__fp_parse_small_leading:wwNN 1 #1 ; #2; #3 #4
17479 {
17480   #1 #2
17481   \exp_after:wN \__fp_parse_pack_trailing:NNNNNNww

```

```

17482 \exp_after:wN 0
17483 \int_value:w \_fp_int_eval:w 1
17484 \if_int_compare:w 9 < 1 \token_to_str:N #4 \exp_stop_f:
17485 \token_to_str:N #4
17486 \exp_after:wN \_fp_parse_small_trailing:wwNN
17487 \int_value:w 1
17488 \exp_after:wN \_fp_parse_digits_vi:N
17489 \exp:w
17490 \else:
17491 0000 0000 \_fp_parse_exponent:Nw #4
17492 \fi:
17493 \_fp_parse_expand:w
17494 }

```

(End definition for _fp_parse_small_leading:wwNN.)

```

\_fp_parse_small_trailing:wwNN \_fp_parse_small_trailing:wwNN 1 <digits> ; <zeros> ; <number of zeros>
<next token>

```

Leave digits 10 to 15 (arguments #1 and #2) in the input stream. If the *<next token>* is a digit, it is the 16th digit, we keep it, then the `small_round` auxiliary considers this digit and all further digits to perform the rounding: the function expands to nothing, to +0 or to +1. Otherwise, there is no 16-th digit, so we put a 0, and look for an exponent.

```

17495 \cs_new:Npn \_fp_parse_small_trailing:wwNN 1 #1 ; #2; #3 #4
17496 {
17497   #1 #2
17498   \if_int_compare:w 9 < 1 \token_to_str:N #4 \exp_stop_f:
17499   \token_to_str:N #4
17500   \exp_after:wN \_fp_parse_small_round:NN
17501   \exp_after:wN #4
17502   \exp:w
17503   \else:
17504   0 \_fp_parse_exponent:Nw #4
17505   \fi:
17506   \_fp_parse_expand:w
17507 }

```

(End definition for _fp_parse_small_trailing:wwNN.)

```

\_fp_parse_pack_trailing:NNNNNNww
\_fp_parse_pack_leading:NNNNNNww
\_fp_parse_pack_carry:w

```

Those functions are expanded after all the digits are found, we took care of the rounding, as well as the exponent. The last argument is the exponent. The previous five arguments are 8 digits which we pack in groups of 4, and the argument before that is 1, except in the rare case where rounding lead to a carry, in which case the argument is 2. The `trailing` function has an exponent shift as its first argument, which we add to the exponent found in the `e...` syntax. If the trailing digits cause a carry, the integer expression for the leading digits is incremented (+1 in the code below). If the leading digits propagate this carry all the way up, the function `_fp_parse_pack_carry:w` increments the exponent, and changes the significand from 0000... to 1000...: this is simple because such a carry can only occur to give rise to a power of 10.

```

17508 \cs_new:Npn \_fp_parse_pack_trailing:NNNNNNww #1 #2 #3#4#5#6 #7; #8 ;
17509 {
17510   \if_meaning:w 2 #2 + 1 \fi:
17511   ; #8 + #1 ; {#3#4#5#6} {#7};
17512 }

```



```

17513 \cs_new:Npn \__fp_parse_pack_leading:NNNNNww #1 #2#3#4#5 #6; #7;
17514 {
17515     + #7
17516     \if_meaning:w 2 #1 \__fp_parse_pack_carry:w \fi:
17517     ; 0 {#2#3#4#5} {#6}
17518 }
17519 \cs_new:Npn \__fp_parse_pack_carry:w \fi: ; 0 #1
17520 { \fi: + 1 ; 0 {1000} }

```

(End definition for __fp_parse_pack_trailing:NNNNNww, __fp_parse_pack_leading:NNNNNww, and __fp_parse_pack_carry:w.)

28.4.3 Number: large significand

Parsing a significand larger than 1 is a little bit more difficult than parsing small significands. We need to count the number of digits before the decimal separator, and add that to the final exponent. We also need to test for the presence of a dot each time we run out of digits, and branch to the appropriate `parse_small` function in those cases.

__fp_parse_large:N This function is followed by the first non-zero digit of a “large” significand (≥ 1). It is called within an integer expression for the exponent. Grab up to 7 more digits, for a total of 8 digits.

```

17521 \cs_new:Npn \__fp_parse_large:N #1
17522 {
17523     \exp_after:wN \__fp_parse_large_leading:wwNN
17524     \int_value:w 1 \token_to_str:N #1
17525     \exp_after:wN \__fp_parse_digits_vii:N
17526     \exp:w \__fp_parse_expand:w
17527 }

```

(End definition for __fp_parse_large:N.)

__fp_parse_large_leading:wwNN __fp_parse_large_leading:wwNN 1 $\langle digits \rangle$; $\langle zeros \rangle$; $\langle number\ of\ zeros \rangle$ $\langle next\ token \rangle$

We shift the exponent by the number of digits in #1, namely the target number, 8, minus the $\langle number\ of\ zeros \rangle$ (number of digits missing). Then prepare to pack the 8 first digits. If the $\langle next\ token \rangle$ is a digit, read up to 6 more digits (digits 10 to 15). If it is a period, try to grab the end of our 8 first digits, branching to the `small` functions since the number of digit does not affect the exponent anymore. Finally, if this is the end of the significand, insert the $\langle zeros \rangle$ to complete the 8 first digits, insert 8 more, and look for an exponent.

```

17528 \cs_new:Npn \__fp_parse_large_leading:wwNN 1 #1 ; #2; #3 #4
17529 {
17530     + \c__fp_half_prec_int - #3
17531     \exp_after:wN \__fp_parse_pack_leading:NNNNNww
17532     \int_value:w \__fp_int_eval:w 1 #1
17533     \if_int_compare:w 9 < 1 \token_to_str:N #4 \exp_stop_f:
17534     \exp_after:wN \__fp_parse_large_trailing:wwNN
17535     \int_value:w 1 \token_to_str:N #4
17536     \exp_after:wN \__fp_parse_digits_vi:N
17537     \exp:w
17538     \else:
17539     \if:w . \exp_not:N #4

```

```

17540         \exp_after:wN \__fp_parse_small_leading:wwNN
17541         \int_value:w 1
17542         \cs:w
17543         __fp_parse_digits_
17544         \__fp_int_to_roman:w #3
17545         :N \exp_after:wN
17546         \cs_end:
17547         \exp:w
17548     \else:
17549         #2
17550         \exp_after:wN \__fp_parse_pack_trailing:NNNNNNww
17551         \exp_after:wN 0
17552         \int_value:w 1 0000 0000
17553         \__fp_parse_exponent:Nw #4
17554     \fi:
17555 \fi:
17556 \__fp_parse_expand:w
17557 }

```

(End definition for __fp_parse_large_leading:wwNN.)

```

\__fp_parse_large_trailing:wwNN
    \__fp_parse_large_trailing:wwNN 1 <digits> ; <zeros> ; <number of zeros>
    <next token>

```

We have just read 15 digits. If the *<next token>* is a digit, then the exponent shift caused by this block of 8 digits is 8, first argument to the `pack_trailing` function. We keep the *<digits>* and this 16-th digit, and find how this should be rounded using `__fp_parse_large_round:NN`. Otherwise, the exponent shift is the number of *<digits>*, 7 minus the *<number of zeros>*, and we test for a decimal point. This case happens in 123451234512345.67 with exactly 15 digits before the decimal separator. Then branch to the appropriate `small` auxiliary, grabbing a few more digits to complement the digits we already grabbed. Finally, if this is truly the end of the significand, look for an exponent after using the *<zeros>* and providing a 16-th digit of 0.

```

17558 \cs_new:Npn \__fp_parse_large_trailing:wwNN 1 #1 ; #2; #3 #4
17559 {
17560     \if_int_compare:w 9 < 1 \token_to_str:N #4 \exp_stop_f:
17561     \exp_after:wN \__fp_parse_pack_trailing:NNNNNNww
17562     \exp_after:wN \c__fp_half_prec_int
17563     \int_value:w \__fp_int_eval:w 1 #1 \token_to_str:N #4
17564     \exp_after:wN \__fp_parse_large_round:NN
17565     \exp_after:wN #4
17566     \exp:w
17567 \else:
17568     \exp_after:wN \__fp_parse_pack_trailing:NNNNNNww
17569     \int_value:w \__fp_int_eval:w 7 - #3 \exp_stop_f:
17570     \int_value:w \__fp_int_eval:w 1 #1
17571     \if:w . \exp_not:N #4
17572     \exp_after:wN \__fp_parse_small_trailing:wwNN
17573     \int_value:w 1
17574     \cs:w
17575     __fp_parse_digits_
17576     \__fp_int_to_roman:w #3
17577     :N \exp_after:wN
17578     \cs_end:

```

```

17579         \exp:w
17580     \else:
17581         #2 0 \__fp_parse_exponent:Nw #4
17582     \fi:
17583 \fi:
17584 \__fp_parse_expand:w
17585 }

```

(End definition for __fp_parse_large_trailing:wwNN.)

28.4.4 Number: beyond 16 digits, rounding

__fp_parse_round_loop:N This loop is called when rounding a number (whether the mantissa is small or large).
 __fp_parse_round_up:N It should appear in an integer expression. This function reads digits one by one, until reaching a non-digit, and adds 1 to the integer expression for each digit. If all digits found are 0, the function ends the expression by ;0, otherwise by ;1. This is done by switching the loop to round_up at the first non-zero digit, thus we avoid to test whether digits are 0 or not once we see a first non-zero digit.

```

17586 \cs_new:Npn \__fp_parse_round_loop:N #1
17587 {
17588     \if_int_compare:w 9 < 1 \token_to_str:N #1 \exp_stop_f:
17589         + 1
17590     \if:w 0 \token_to_str:N #1
17591         \exp_after:wN \__fp_parse_round_loop:N
17592         \exp:w
17593     \else:
17594         \exp_after:wN \__fp_parse_round_up:N
17595         \exp:w
17596     \fi:
17597 \else:
17598     \__fp_parse_return_semicolon:w 0 #1
17599 \fi:
17600 \__fp_parse_expand:w
17601 }
17602 \cs_new:Npn \__fp_parse_round_up:N #1
17603 {
17604     \if_int_compare:w 9 < 1 \token_to_str:N #1 \exp_stop_f:
17605         + 1
17606     \exp_after:wN \__fp_parse_round_up:N
17607     \exp:w
17608 \else:
17609     \__fp_parse_return_semicolon:w 1 #1
17610 \fi:
17611 \__fp_parse_expand:w
17612 }

```

(End definition for __fp_parse_round_loop:N and __fp_parse_round_up:N.)

__fp_parse_round_after:wN After the loop __fp_parse_round_loop:N, this function fetches an exponent with __fp_parse_exponent:N, and combines it with the number of digits counted by __fp_parse_round_loop:N. At the same time, the result 0 or 1 is added to the surrounding integer expression.

```

17613 \cs_new:Npn \__fp_parse_round_after:wN #1; #2
17614 {

```

```

17615     + #2 \exp_after:wN ;
17616     \int_value:w \__fp_int_eval:w #1 + \__fp_parse_exponent:N
17617 }

```

(End definition for __fp_parse_round_after:wN.)

__fp_parse_small_round:NN Here, #1 is the digit that we are currently rounding (we only care whether it is even or odd). If #2 is not a digit, then fetch an exponent and expand to ;*<exponent>* only. Otherwise, we expand to +0 or +1, then ;*<exponent>*. To decide which, call __fp_round_s:NNNw to know whether to round up, giving it as arguments a sign 0 (all explicit numbers are positive), the digit #1 to round, the first following digit #2, and either +0 or +1 depending on whether the following digits are all zero or not. This last argument is obtained by __fp_parse_round_loop:N, whose number of digits we discard by multiplying it by 0. The exponent which follows the number is also fetched by __fp_parse_round_after:wN.

```

17618 \cs_new:Npn \__fp_parse_small_round:NN #1#2
17619 {
17620   \if_int_compare:w 9 < 1 \token_to_str:N #2 \exp_stop_f:
17621   +
17622   \exp_after:wN \__fp_round_s:NNNw
17623   \exp_after:wN 0
17624   \exp_after:wN #1
17625   \exp_after:wN #2
17626   \int_value:w \__fp_int_eval:w
17627   \exp_after:wN \__fp_parse_round_after:wN
17628   \int_value:w \__fp_int_eval:w 0 * \__fp_int_eval:w 0
17629   \exp_after:wN \__fp_parse_round_loop:N
17630   \exp:w
17631   \else:
17632     \__fp_parse_exponent:Nw #2
17633   \fi:
17634   \__fp_parse_expand:w
17635 }

```

(End definition for __fp_parse_small_round:NN and __fp_parse_round_after:wN.)

__fp_parse_large_round:NN Large numbers are harder to round, as there may be a period in the way. Again, #1 is the digit that we are currently rounding (we only care whether it is even or odd). If there are no more digits (#2 is not a digit), then we must test for a period: if there is one, then switch to the rounding function for small significands, otherwise fetch an exponent. If there are more digits (#2 is a digit), then round, checking with __fp_parse_round_loop:N if all further digits vanish, or some are non-zero. This loop is not enough, as it is stopped by a period. After the loop, the aux function tests for a period: if it is present, then we must continue looking for digits, this time discarding the number of digits we find.

```

17636 \cs_new:Npn \__fp_parse_large_round:NN #1#2
17637 {
17638   \if_int_compare:w 9 < 1 \token_to_str:N #2 \exp_stop_f:
17639   +
17640   \exp_after:wN \__fp_round_s:NNNw
17641   \exp_after:wN 0
17642   \exp_after:wN #1
17643   \exp_after:wN #2

```

```

17644     \int_value:w \__fp_int_eval:w
17645     \exp_after:wN \__fp_parse_large_round_aux:wNN
17646     \int_value:w \__fp_int_eval:w 1
17647     \exp_after:wN \__fp_parse_round_loop:N
17648 \else: %^^A could be dot, or e, or other
17649     \exp_after:wN \__fp_parse_large_round_test:NN
17650     \exp_after:wN #1
17651     \exp_after:wN #2
17652 \fi:
17653 }
17654 \cs_new:Npn \__fp_parse_large_round_test:NN #1#2
17655 {
17656     \if:w . \exp_not:N #2
17657         \exp_after:wN \__fp_parse_small_round:NN
17658         \exp_after:wN #1
17659         \exp:w
17660     \else:
17661         \__fp_parse_exponent:Nw #2
17662     \fi:
17663     \__fp_parse_expand:w
17664 }
17665 \cs_new:Npn \__fp_parse_large_round_aux:wNN #1 ; #2 #3
17666 {
17667     + #2
17668     \exp_after:wN \__fp_parse_round_after:wN
17669     \int_value:w \__fp_int_eval:w #1
17670     \if:w . \exp_not:N #3
17671         + 0 * \__fp_int_eval:w 0
17672         \exp_after:wN \__fp_parse_round_loop:N
17673         \exp:w \exp_after:wN \__fp_parse_expand:w
17674     \else:
17675         \exp_after:wN ;
17676         \exp_after:wN 0
17677         \exp_after:wN #3
17678     \fi:
17679 }

```

(End definition for `__fp_parse_large_round:NN`, `__fp_parse_large_round_test:NN`, and `__fp_parse_large_round_aux:wNN`.)

28.4.5 Number: finding the exponent

Expansion is a little bit tricky here, in part because we accept input where multiplication is implicit.

```

\__fp_parse:n { 3.2 erf(0.1) }
\__fp_parse:n { 3.2 e\l_my_int }
\__fp_parse:n { 3.2 \c_pi_fp }

```

The first case indicates that just looking one character ahead for an “e” is not enough, since we would mistake the function `erf` for an exponent of “rf”. An alternative would be to look two tokens ahead and check if what follows is a sign or a digit, considering in that case that we must be finding an exponent. But taking care of the second case requires that we unpack registers after `e`. However, blindly expanding the two tokens ahead completely would break the third example (unpacking is even worse). Indeed, in

the course of reading 3.2, `\c_pi_fp` is expanded to `\s__fp __fp_chk:w 1 0 {-1} {3141}` `__fp` stops the expansion. Expanding two tokens ahead would then force the expansion of `__fp_chk:w` (despite it being protected), and that function tries to produce an error.

What can we do? Really, the reason why this last case breaks is that just as `TeX` does, we should read ahead as little as possible. Here, the only case where there may be an exponent is if the first token ahead is `e`. Then we expand (and possibly unpack) the second token.

`__fp_parse_exponent:Nw` This auxiliary is convenient to smuggle some material through `\fi:` ending conditional processing. We place those `\fi:` (argument #2) at a very odd place because this allows us to insert `__fp_int_eval:w ...` there if needed.

```
17680 \cs_new:Npn \__fp_parse_exponent:Nw #1 #2 \__fp_parse_expand:w
17681 {
17682   \exp_after:wN ;
17683   \int_value:w #2 \__fp_parse_exponent:N #1
17684 }
```

(End definition for `__fp_parse_exponent:Nw`.)

`__fp_parse_exponent:N`
`__fp_parse_exponent_aux:NN` This function should be called within an `\int_value:w` expansion (or within an integer expression). It leaves digits of the exponent behind it in the input stream, and terminates the expansion with a semicolon. If there is no `e` (or `E`), leave an exponent of 0. If there is an `e` or `E`, expand the next token to run some tests on it. The first rough test is that if the character code of #1 is greater than that of 9 (largest code valid for an exponent, less than any code valid for an identifier), there was in fact no exponent; otherwise, we search for the sign of the exponent.

```
17685 \cs_new:Npn \__fp_parse_exponent:N #1
17686 {
17687   \if:w e \if:w E \exp_not:N #1 e \else: \exp_not:N #1 \fi:
17688   \exp_after:wN \__fp_parse_exponent_aux:NN
17689   \exp_after:wN #1
17690   \exp:w
17691   \else:
17692     0 \__fp_parse_return_semicolon:w #1
17693   \fi:
17694   \__fp_parse_expand:w
17695 }
17696 \cs_new:Npn \__fp_parse_exponent_aux:NN #1#2
17697 {
17698   \if_int_compare:w \if_catcode:w \scan_stop: \exp_not:N #2
17699     0 \else: '#2 \fi: > '9 \exp_stop_f:
17700     0 \exp_after:wN ; \exp_after:wN #1
17701   \else:
17702     \exp_after:wN \__fp_parse_exponent_sign:N
17703   \fi:
17704   #2
17705 }
```

(End definition for `__fp_parse_exponent:N` and `__fp_parse_exponent_aux:NN`.)

`__fp_parse_exponent_sign:N` Read signs one by one (if there is any).

```
17706 \cs_new:Npn \__fp_parse_exponent_sign:N #1
```

```

17707 {
17708   \if:w + \if:w - \exp_not:N #1 + \fi: \token_to_str:N #1
17709   \exp_after:wN \__fp_parse_exponent_sign:N
17710   \exp:w \exp_after:wN \__fp_parse_expand:w
17711   \else:
17712     \exp_after:wN \__fp_parse_exponent_body:N
17713     \exp_after:wN #1
17714   \fi:
17715 }

```

(End definition for `__fp_parse_exponent_sign:N`.)

`__fp_parse_exponent_body:N` An exponent can be an explicit integer (most common case), or various other things (most of which are invalid).

```

17716 \cs_new:Npn \__fp_parse_exponent_body:N #1
17717 {
17718   \if_int_compare:w 9 < 1 \token_to_str:N #1 \exp_stop_f:
17719   \token_to_str:N #1
17720   \exp_after:wN \__fp_parse_exponent_digits:N
17721   \exp:w
17722   \else:
17723     \__fp_parse_exponent_keep:NTF #1
17724     { \__fp_parse_return_semicolon:w #1 }
17725     {
17726       \exp_after:wN ;
17727       \exp:w
17728     }
17729   \fi:
17730   \__fp_parse_expand:w
17731 }

```

(End definition for `__fp_parse_exponent_body:N`.)

`__fp_parse_exponent_digits:N` Read digits one by one, and leave them behind in the input stream. When finding a non-digit, stop, and insert a semicolon. Note that we do not check for overflow of the exponent, hence there can be a T_EX error. It is mostly harmless, except when parsing 0e9876543210, which should be a valid representation of 0, but is not.

```

17732 \cs_new:Npn \__fp_parse_exponent_digits:N #1
17733 {
17734   \if_int_compare:w 9 < 1 \token_to_str:N #1 \exp_stop_f:
17735   \token_to_str:N #1
17736   \exp_after:wN \__fp_parse_exponent_digits:N
17737   \exp:w
17738   \else:
17739     \__fp_parse_return_semicolon:w #1
17740   \fi:
17741   \__fp_parse_expand:w
17742 }

```

(End definition for `__fp_parse_exponent_digits:N`.)

`__fp_parse_exponent_keep:NTF` This is the last building block for parsing exponents. The argument `#1` is already fully expanded, and neither `+` nor `-` nor a digit. It can be:

- `\s__fp`, marking the start of an internal floating point, invalid here;

- another control sequence equal to `\relax`, probably a bad variable;
- a register: in this case we make sure that it is an integer register, not a dimension;
- a character other than +, - or digits, again, an error.

```

17743 \prg_new_conditional:Npnn \_fp_parse_exponent_keep:N #1 { TF }
17744 {
17745   \if_catcode:w \scan_stop: \exp_not:N #1
17746   \if_meaning:w \scan_stop: #1
17747   \if_int_compare:w
17748     \_fp_str_if_eq:nn { \s_fp } { \exp_not:N #1 }
17749     = 0 \exp_stop_f:
17750     0
17751     \_kernel_msg_expandable_error:nnn
17752     { kernel } { fp-after-e } { floating-point~ }
17753     \prg_return_true:
17754   \else:
17755     0
17756     \_kernel_msg_expandable_error:nnn
17757     { kernel } { bad-variable } { #1 }
17758     \prg_return_false:
17759   \fi:
17760 \else:
17761   \if_int_compare:w
17762     \_fp_str_if_eq:nn { \int_value:w #1 } { \tex_the:D #1 }
17763     = 0 \exp_stop_f:
17764     \int_value:w #1
17765   \else:
17766     0
17767     \_kernel_msg_expandable_error:nnn
17768     { kernel } { fp-after-e } { dimension~#1 }
17769   \fi:
17770   \prg_return_false:
17771 \fi:
17772 \else:
17773   0
17774   \_kernel_msg_expandable_error:nnn
17775   { kernel } { fp-missing } { exponent }
17776   \prg_return_true:
17777 \fi:
17778 }

```

(End definition for `_fp_parse_exponent_keep:N`.)

28.5 Constants, functions and prefix operators

28.5.1 Prefix operators

`_fp_parse_prefix_+:Nw` A unary + does nothing: we should continue looking for a number.

```

17779 \cs_new_eq:cN { \_fp_parse_prefix_+:Nw } \_fp_parse_one:Nw

```

(End definition for `_fp_parse_prefix_+:Nw`.)

_fp_parse_apply_function:NNwN

Here, #1 is a precedence, #2 is some extra data used by some functions, #3 is *e.g.*, _fp_sin_o:w, and expands once after the calculation, #4 is the operand, and #5 is a _fp_parse_infix...:N function. We feed the data #2, and the argument #4, to the function #3, which expands \exp:w thus the infix function #5.

```
17780 \cs_new:Npn \_fp_parse_apply_function:NNwN #1#2#3#4#5
17781 {
17782     #3 #2 #4 @
17783     \exp:w \exp_end_continue_f:w #5 #1
17784 }
```

(End definition for _fp_parse_apply_function:NNwN.)

_fp_parse_apply_unary:NNwN

_fp_parse_apply_unary_chk:NwNw

_fp_parse_apply_unary_chk:nNNNw

_fp_parse_apply_unary_type:NNN

_fp_parse_apply_unary_error:NNw

In contrast to _fp_parse_apply_function:NNwN, this checks that the operand #4 is a single argument (namely there is a single ;). We use the fact that any floating point starts with a “safe” token like \s__fp. If there is no argument produce the fp-no-arg error; if there are at least two produce fp-multi-arg. For the error message extract the mathematical function name (such as sin) from the expl3 function that computes it, such as _fp_sin_o:w.

In addition, since there is a single argument we can dispatch on type and check that the resulting function exists. This catches things like sin((1,2)) where it does not make sense to take the sine of a tuple.

```
17785 \cs_new:Npn \_fp_parse_apply_unary:NNwN #1#2#3#4#5
17786 {
17787     \_fp_parse_apply_unary_chk:NwNw #4 @ ; . \s__fp_stop
17788     \_fp_parse_apply_unary_type:NNN
17789     #3 #2 #4 @
17790     \exp:w \exp_end_continue_f:w #5 #1
17791 }
17792 \cs_new:Npn \_fp_parse_apply_unary_chk:NwNw #1#2 ; #3#4 \s__fp_stop
17793 {
17794     \if_meaning:w @ #3 \else:
17795         \token_if_eq_meaning:NNTF . #3
17796         { \_fp_parse_apply_unary_chk:nNNNNw { no } }
17797         { \_fp_parse_apply_unary_chk:nNNNNw { multi } }
17798     \fi:
17799 }
17800 \cs_new:Npn \_fp_parse_apply_unary_chk:nNNNNw #1#2#3#4#5#6 @
17801 {
17802     #2
17803     \_fp_error:nffn { fp-#1-arg } { \_fp_func_to_name:N #4 } { } { }
17804     \exp_after:wN #4 \exp_after:wN #5 \c_nan_fp @
17805 }
17806 \cs_new:Npn \_fp_parse_apply_unary_type:NNN #1#2#3
17807 {
17808     \_fp_change_func_type:NNN #3 #1 \_fp_parse_apply_unary_error:NNw
17809     #2 #3
17810 }
17811 \cs_new:Npn \_fp_parse_apply_unary_error:NNw #1#2#3 @
17812 { \_fp_invalid_operation_o:fw { \_fp_func_to_name:N #1 } #3 }
```

(End definition for _fp_parse_apply_unary:NNwN and others.)

`__fp_parse_prefix_-:Nw` The unary `-` and boolean not are harder: we parse the operand using a precedence equal
`__fp_parse_prefix_!:Nw` to the maximum of the previous precedence `##1` and the precedence `\c__fp_prec_not_-int` of the unary operator, then call the appropriate `__fp_⟨operation⟩_o:w` function, where the `⟨operation⟩` is `set_sign` or `not`.

```

17813 \cs_set_protected:Npn \__fp_tmp:w #1#2#3#4
17814 {
17815   \cs_new:cpn { __fp_parse_prefix_ #1 :Nw } ##1
17816   {
17817     \exp_after:wN \__fp_parse_apply_unary:NNwN
17818     \exp_after:wN ##1
17819     \exp_after:wN #4
17820     \exp_after:wN #3
17821     \exp:w
17822     \if_int_compare:w #2 < ##1
17823       \__fp_parse_operand:Nw ##1
17824     \else:
17825       \__fp_parse_operand:Nw #2
17826     \fi:
17827     \__fp_parse_expand:w
17828   }
17829 }
17830 \__fp_tmp:w - \c__fp_prec_not_int \__fp_set_sign_o:w 2
17831 \__fp_tmp:w ! \c__fp_prec_not_int \__fp_not_o:w ?

```

(End definition for `__fp_parse_prefix_-:Nw` and `__fp_parse_prefix_!:Nw`.)

`__fp_parse_prefix_:Nw` Numbers which start with a decimal separator (a period) end up here. Of course, we do not look for an operand, but for the rest of the number. This function is very similar to `__fp_parse_one_digit:NN` but calls `__fp_parse_strim_zeros:N` to trim zeros after the decimal point, rather than the `trim_zeros` function for zeros before the decimal point.

```

17832 \cs_new:cpn { __fp_parse_prefix_:Nw } #1
17833 {
17834   \exp_after:wN \__fp_parse_infix_after_operand:NwN
17835   \exp_after:wN #1
17836   \exp:w \exp_end_continue_f:w
17837   \exp_after:wN \__fp_sanitize:wN
17838   \int_value:w \__fp_int_eval:w 0 \__fp_parse_strim_zeros:N
17839 }

```

(End definition for `__fp_parse_prefix_:Nw`.)

`__fp_parse_prefix_(:Nw` The left parenthesis is treated as a unary prefix operator because it appears in exactly
`__fp_parse_lparen_after:NwN` the same settings. If the previous precedence is `\c__fp_prec_func_int` we are parsing arguments of a function and commas should not build tuples; otherwise commas should build tuples. We distinguish these cases by precedence: `\c__fp_prec_comma_int` for the case of arguments, `\c__fp_prec_tuple_int` for the case of tuples. Once the operand is found, the `lparen_after` auxiliary makes sure that there was a closing parenthesis (otherwise it complains), and leaves in the input stream an operand, fetching the following infix operator.

```

17840 \cs_new:cpn { __fp_parse_prefix_(:Nw } #1
17841 {
17842   \exp_after:wN \__fp_parse_lparen_after:NwN

```

```

17843 \exp_after:wN #1
17844 \exp:w
17845 \if_int_compare:w #1 = \c__fp_prec_func_int
17846 \__fp_parse_operand:Nw \c__fp_prec_comma_int
17847 \else:
17848 \__fp_parse_operand:Nw \c__fp_prec_tuple_int
17849 \fi:
17850 \__fp_parse_expand:w
17851 }
17852 \cs_new:Npx \__fp_parse_lparen_after:NwN #1#2 @ #3
17853 {
17854 \exp_not:N \token_if_eq_meaning:NNTF #3
17855 \exp_not:c { __fp_parse_infix_ }:N }
17856 {
17857 \exp_not:N \__fp_exp_after_array_f:w #2 \s__fp_expr_stop
17858 \exp_not:N \exp_after:wN
17859 \exp_not:N \__fp_parse_infix_after_paren:NN
17860 \exp_not:N \exp_after:wN #1
17861 \exp_not:N \exp:w
17862 \exp_not:N \__fp_parse_expand:w
17863 }
17864 {
17865 \exp_not:N \__kernel_msg_expandable_error:nnn
17866 { kernel } { fp-missing } { ) }
17867 \exp_not:N \tl_if_empty:nT {#2} \exp_not:N \c__fp_empty_tuple_fp
17868 #2 @
17869 \exp_not:N \use_none:n #3
17870 }
17871 }

```

(End definition for __fp_parse_prefix_(:Nw and __fp_parse_lparen_after:NwN.)

__fp_parse_prefix_):Nw The right parenthesis can appear as a prefix in two similar cases: in an empty tuple or tuple ending with a comma, or in an empty argument list or argument list ending with a comma, such as in `max(1,2,)` or in `rand()`.

```

17872 \cs_new:cpn { __fp_parse_prefix_):Nw } #1
17873 {
17874 \if_int_compare:w #1 = \c__fp_prec_comma_int
17875 \else:
17876 \if_int_compare:w #1 = \c__fp_prec_tuple_int
17877 \exp_after:wN \c__fp_empty_tuple_fp \exp:w
17878 \else:
17879 \__kernel_msg_expandable_error:nnn
17880 { kernel } { fp-missing-number } { ) }
17881 \exp_after:wN \c_nan_fp \exp:w
17882 \fi:
17883 \exp_end_continue_f:w
17884 \fi:
17885 \__fp_parse_infix_after_paren:NN #1 )
17886 }

```

(End definition for __fp_parse_prefix_):Nw.)

28.5.2 Constants

Some words correspond to constant floating points. The floating point constant is left as a result of `__fp_parse_one:Nw` after expanding `__fp_parse_infix:NN`.

```

\__fp_parse_word_inf:N
\__fp_parse_word_nan:N
\__fp_parse_word_pi:N
\__fp_parse_word_deg:N
\__fp_parse_word_true:N
\__fp_parse_word_false:N
17887 \cs_set_protected:Npn \__fp_tmp:w #1 #2
17888 {
17889   \cs_new:cpn { __fp_parse_word_#1:N }
17890     { \exp_after:wN #2 \exp:w \exp_end_continue_f:w \__fp_parse_infix:NN }
17891 }

```

```

17892 \__fp_tmp:w { inf } \c_inf_fp
17893 \__fp_tmp:w { nan } \c_nan_fp
17894 \__fp_tmp:w { pi } \c_pi_fp
17895 \__fp_tmp:w { deg } \c_one_degree_fp
17896 \__fp_tmp:w { true } \c_one_fp
17897 \__fp_tmp:w { false } \c_zero_fp

```

(End definition for `__fp_parse_word_inf:N` and others.)

Copies of `__fp_parse_word_...:N` commands, to allow arbitrary case as mandated by the standard.

```

\__fp_parse_caseless_inf:N
\__fp_parse_caseless_infinity:N
\__fp_parse_caseless_nan:N
17898 \cs_new_eq:NN \__fp_parse_caseless_inf:N \__fp_parse_word_inf:N
17899 \cs_new_eq:NN \__fp_parse_caseless_infinity:N \__fp_parse_word_inf:N
17900 \cs_new_eq:NN \__fp_parse_caseless_nan:N \__fp_parse_word_nan:N

```

(End definition for `__fp_parse_caseless_inf:N`, `__fp_parse_caseless_infinity:N`, and `__fp_parse_caseless_nan:N`.)

Dimension units are also floating point constants but their value is not stored as a floating point constant. We give the values explicitly here.

```

\__fp_parse_word_pt:N
\__fp_parse_word_in:N
\__fp_parse_word_pc:N
\__fp_parse_word_cm:N
\__fp_parse_word_mm:N
\__fp_parse_word_dd:N
\__fp_parse_word_cc:N
\__fp_parse_word_nd:N
\__fp_parse_word_nc:N
\__fp_parse_word_bp:N
\__fp_parse_word_sp:N
17901 \cs_set_protected:Npn \__fp_tmp:w #1 #2
17902 {
17903   \cs_new:cpn { __fp_parse_word_#1:N }
17904     {
17905       \__fp_exp_after_f:nw { \__fp_parse_infix:NN }
17906       \s__fp \__fp_chk:w 10 #2 ;
17907     }
17908 }
17909 \__fp_tmp:w {pt} { {1} {1000} {0000} {0000} {0000} }
17910 \__fp_tmp:w {in} { {2} {7227} {0000} {0000} {0000} }
17911 \__fp_tmp:w {pc} { {2} {1200} {0000} {0000} {0000} }
17912 \__fp_tmp:w {cm} { {2} {2845} {2755} {9055} {1181} }
17913 \__fp_tmp:w {mm} { {1} {2845} {2755} {9055} {1181} }
17914 \__fp_tmp:w {dd} { {1} {1070} {0085} {6496} {0630} }
17915 \__fp_tmp:w {cc} { {2} {1284} {0102} {7795} {2756} }
17916 \__fp_tmp:w {nd} { {1} {1066} {9783} {4645} {6693} }
17917 \__fp_tmp:w {nc} { {2} {1280} {3740} {1574} {8031} }
17918 \__fp_tmp:w {bp} { {1} {1003} {7500} {0000} {0000} }
17919 \__fp_tmp:w {sp} { {-4} {1525} {8789} {0625} {0000} }

```

(End definition for `__fp_parse_word_pt:N` and others.)

The font-dependent units `em` and `ex` must be evaluated on the fly. We reuse an auxiliary of `\dim_to_fp:n`.

```

\__fp_parse_word_em:N
\__fp_parse_word_ex:N
17920 \tl_map_inline:nn { {em} {ex} }
17921 {

```

```

17922 \cs_new:cpn { __fp_parse_word_#1:N }
17923 {
17924   \exp_after:wN \__fp_from_dim_test:ww
17925   \exp_after:wN 0 \exp_after:wN ,
17926   \int_value:w \dim_to_decimal_in_sp:n { 1 #1 } \exp_after:wN ;
17927   \exp:w \exp_end_continue_f:w \__fp_parse_infix:NN
17928 }
17929 }

```

(End definition for `__fp_parse_word_em:N` and `__fp_parse_word_ex:N`.)

28.5.3 Functions

```

\__fp_parse_unary_function:NNN
\__fp_parse_function:NNN
17930 \cs_new:Npn \__fp_parse_unary_function:NNN #1#2#3
17931 {
17932   \exp_after:wN \__fp_parse_apply_unary:NNNwN
17933   \exp_after:wN #3
17934   \exp_after:wN #2
17935   \exp_after:wN #1
17936   \exp:w
17937   \__fp_parse_operand:Nw \c__fp_prec_func_int \__fp_parse_expand:w
17938 }
17939 \cs_new:Npn \__fp_parse_function:NNN #1#2#3
17940 {
17941   \exp_after:wN \__fp_parse_apply_function:NNNwN
17942   \exp_after:wN #3
17943   \exp_after:wN #2
17944   \exp_after:wN #1
17945   \exp:w
17946   \__fp_parse_operand:Nw \c__fp_prec_func_int \__fp_parse_expand:w
17947 }

```

(End definition for `__fp_parse_unary_function:NNN` and `__fp_parse_function:NNN`.)

28.6 Main functions

`__fp_parse:n` Start an `\exp:w` expansion so that `__fp_parse:n` expands in two steps. The `__fp_parse_operand:Nw` function performs computations until reaching an operation with precedence `\c__fp_prec_end_int` or less, namely, the end of the expression. The marker `\s__fp_expr_mark` indicates that the next token is an already parsed version of an infix operator, and `__fp_parse_infix_end:N` has infinitely negative precedence. Finally, clean up a (well-defined) set of extra tokens and stop the initial expansion with `\exp_end:.`

```

17948 \cs_new:Npn \__fp_parse:n #1
17949 {
17950   \exp:w
17951   \exp_after:wN \__fp_parse_after:ww
17952   \exp:w
17953   \__fp_parse_operand:Nw \c__fp_prec_end_int
17954   \__fp_parse_expand:w #1
17955   \s__fp_expr_mark \__fp_parse_infix_end:N
17956   \s__fp_expr_stop

```

```

17957     \exp_end:
17958   }
17959   \cs_new:Npn \__fp_parse_after:ww
17960     #1@ \__fp_parse_infix_end:N \s__fp_expr_stop #2 { #2 #1 }
17961   \cs_new:Npn \__fp_parse_o:n #1
17962   {
17963     \exp:w
17964     \exp_after:wN \__fp_parse_after:ww
17965     \exp:w
17966     \__fp_parse_operand:Nw \c__fp_prec_end_int
17967     \__fp_parse_expand:w #1
17968     \s__fp_expr_mark \__fp_parse_infix_end:N
17969     \s__fp_expr_stop
17970   {
17971     \exp_end_continue_f:w
17972     \__fp_exp_after_any_f:nw { \exp_after:wN \exp_stop_f: }
17973   }
17974 }

```

(End definition for __fp_parse:n, __fp_parse_o:n, and __fp_parse_after:ww.)

__fp_parse_operand:Nw This is just a shorthand which sets up both __fp_parse_continue:NwN and __fp_parse_one:Nw with the same precedence. Note the trailing \exp:w.

```

17975 \cs_new:Npn \__fp_parse_operand:Nw #1
17976 {
17977   \exp_end_continue_f:w
17978   \exp_after:wN \__fp_parse_continue:NwN
17979   \exp_after:wN #1
17980   \exp:w \exp_end_continue_f:w
17981   \exp_after:wN \__fp_parse_one:Nw
17982   \exp_after:wN #1
17983   \exp:w
17984 }
17985 \cs_new:Npn \__fp_parse_continue:NwN #1 #2 @ #3 { #3 #1 #2 @ }

```

(End definition for __fp_parse_operand:Nw and __fp_parse_continue:NwN.)

__fp_parse_apply_binary:NwNwN Receives $\langle precedence \rangle \langle operand_1 \rangle @ \langle operation \rangle \langle operand_2 \rangle @ \langle infix command \rangle$. Builds the appropriate call to the $\langle operation \rangle$ #3, dispatching on both types. If the resulting control sequence does not exist, the operation is not allowed.

This is redefined in l3fp-extras.

```

17986 \cs_new:Npn \__fp_parse_apply_binary:NwNwN #1 #2#3@ #4 #5#6@ #7
17987 {
17988   \exp_after:wN \__fp_parse_continue:NwN
17989   \exp_after:wN #1
17990   \exp:w \exp_end_continue_f:w
17991   \exp_after:wN \__fp_parse_apply_binary_chk:NN
17992   \cs:w
17993   __fp
17994   \__fp_type_from_scan:N #2
17995   _#4
17996   \__fp_type_from_scan:N #5
17997   _o:ww
17998   \cs_end:

```

```

17999         #4
18000         #2#3 #5#6
18001         \exp:w \exp_end_continue_f:w #7 #1
18002     }
18003 \cs_new:Npn \__fp_parse_apply_binary_chk:NN #1#2
18004 {
18005     \if_meaning:w \scan_stop: #1
18006     \__fp_parse_apply_binary_error:NNN #2
18007     \fi:
18008     #1
18009 }
18010 \cs_new:Npn \__fp_parse_apply_binary_error:NNN #1#2#3
18011 {
18012     #2
18013     \__fp_invalid_operation_o:Nww #1
18014 }

```

(End definition for __fp_parse_apply_binary:NwNwN, __fp_parse_apply_binary_chk:NN, and __fp_parse_apply_binary_error:NNN.)

__fp_binary_type_o:Nww
__fp_binary_rev_type_o:Nww

Applies the operator #1 to its two arguments, dispatching according to their types, and expands once after the result. The rev version swaps its arguments before doing this.

```

18015 \cs_new:Npn \__fp_binary_type_o:Nww #1 #2#3 ; #4
18016 {
18017     \exp_after:wN \__fp_parse_apply_binary_chk:NN
18018     \cs:w
18019     __fp
18020     \__fp_type_from_scan:N #2
18021     _ #1
18022     \__fp_type_from_scan:N #4
18023     _o:ww
18024     \cs_end:
18025     #1
18026     #2 #3 ; #4
18027 }
18028 \cs_new:Npn \__fp_binary_rev_type_o:Nww #1 #2#3 ; #4#5 ;
18029 {
18030     \exp_after:wN \__fp_parse_apply_binary_chk:NN
18031     \cs:w
18032     __fp
18033     \__fp_type_from_scan:N #4
18034     _ #1
18035     \__fp_type_from_scan:N #2
18036     _o:ww
18037     \cs_end:
18038     #1
18039     #4 #5 ; #2 #3 ;
18040 }

```

(End definition for __fp_binary_type_o:Nww and __fp_binary_rev_type_o:Nww.)

28.7 Infix operators

__fp_parse_infix_after_operand:NwN

```

18041 \cs_new:Npn \__fp_parse_infix_after_operand:NwN #1 #2;
18042 {
18043   \__fp_exp_after_f:nw { \__fp_parse_infix:NN #1 }
18044   #2;
18045 }
18046 \cs_new:Npn \__fp_parse_infix:NN #1 #2
18047 {
18048   \if_catcode:w \scan_stop: \exp_not:N #2
18049   \if_int_compare:w
18050     \__fp_str_if_eq:nn { \s__fp_expr_mark } { \exp_not:N #2 }
18051     = 0 \exp_stop_f:
18052     \exp_after:wN \exp_after:wN
18053     \exp_after:wN \__fp_parse_infix_mark:NNN
18054   \else:
18055     \exp_after:wN \exp_after:wN
18056     \exp_after:wN \__fp_parse_infix_juxt:N
18057   \fi:
18058 \else:
18059   \if_int_compare:w
18060     \__fp_int_eval:w
18061     ( '#2 \if_int_compare:w '#2 > 'Z - 32 \fi: ) / 26
18062     = 3 \exp_stop_f:
18063     \exp_after:wN \exp_after:wN
18064     \exp_after:wN \__fp_parse_infix_juxt:N
18065   \else:
18066     \exp_after:wN \__fp_parse_infix_check:NNN
18067     \cs:w
18068       __fp_parse_infix_ \token_to_str:N #2 :N
18069     \exp_after:wN \exp_after:wN \exp_after:wN
18070     \cs_end:
18071   \fi:
18072 \fi:
18073 #1
18074 #2
18075 }
18076 \cs_new:Npn \__fp_parse_infix_check:NNN #1#2#3
18077 {
18078   \if_meaning:w \scan_stop: #1
18079   \__kernel_msg_expandable_error:nnn
18080   { kernel } { fp-missing } { * }
18081   \exp_after:wN \__fp_parse_infix_mul:N
18082   \exp_after:wN #2
18083   \exp_after:wN #3
18084 \else:
18085   \exp_after:wN #1
18086   \exp_after:wN #2
18087   \exp:w \exp_after:wN \__fp_parse_expand:w
18088 \fi:
18089 }

```

(End definition for __fp_parse_infix_after_operand:NwN.)

__fp_parse_infix_after_paren:NN Variant of __fp_parse_infix:NN for use after a closing parenthesis. The only difference is that __fp_parse_infix_juxt:N is replaced by __fp_parse_infix_mul:N.


```

18090 \cs_new:Npn \__fp_parse_infix_after_paren:NN #1 #2
18091 {
18092   \if_catcode:w \scan_stop: \exp_not:N #2
18093   \if_int_compare:w
18094     \__fp_str_if_eq:nn { \s__fp_expr_mark } { \exp_not:N #2 }
18095     = 0 \exp_stop_f:
18096     \exp_after:wN \exp_after:wN
18097     \exp_after:wN \__fp_parse_infix_mark:NNN
18098   \else:
18099     \exp_after:wN \exp_after:wN
18100     \exp_after:wN \__fp_parse_infix_mul:N
18101   \fi:
18102 \else:
18103   \if_int_compare:w
18104     \__fp_int_eval:w
18105     ( '#2 \if_int_compare:w '#2 > 'Z - 32 \fi: ) / 26
18106     = 3 \exp_stop_f:
18107     \exp_after:wN \exp_after:wN
18108     \exp_after:wN \__fp_parse_infix_mul:N
18109   \else:
18110     \exp_after:wN \__fp_parse_infix_check:NNN
18111     \cs:w
18112       __fp_parse_infix_ \token_to_str:N #2 :N
18113     \exp_after:wN \exp_after:wN \exp_after:wN
18114     \cs_end:
18115   \fi:
18116 \fi:
18117 #1
18118 #2
18119 }

```

(End definition for __fp_parse_infix_after_paren:NN.)

28.7.1 Closing parentheses and commas

__fp_parse_infix_mark:NNN As an infix operator, \s__fp_expr_mark means that the next token (#3) has already gone through __fp_parse_infix:NN and should be provided the precedence #1. The scan mark #2 is discarded.

```

18120 \cs_new:Npn \__fp_parse_infix_mark:NNN #1#2#3 { #3 #1 }

```

(End definition for __fp_parse_infix_mark:NNN.)

__fp_parse_infix_end:N This one is a little bit odd: force every previous operator to end, regardless of the precedence.

```

18121 \cs_new:Npn \__fp_parse_infix_end:N #1
18122 { @ \use_none:n \__fp_parse_infix_end:N }

```

(End definition for __fp_parse_infix_end:N.)

__fp_parse_infix_):N This is very similar to __fp_parse_infix_end:N, complaining about an extra closing parenthesis if the previous operator was the beginning of the expression, with precedence \c__fp_prec_end_int.

```

18123 \cs_set_protected:Npn \__fp_tmp:w #1
18124 {

```

```

18125 \cs_new:Npn #1 ##1
18126 {
18127   \if_int_compare:w ##1 > \c__fp_prec_end_int
18128     \exp_after:wN @
18129     \exp_after:wN \use_none:n
18130     \exp_after:wN #1
18131   \else:
18132     \__kernel_msg_expandable_error:nnn { kernel } { fp-extra } { ) }
18133     \exp_after:wN \__fp_parse_infix:NN
18134     \exp_after:wN ##1
18135     \exp:w \exp_after:wN \__fp_parse_expand:w
18136   \fi:
18137 }
18138 }
18139 \exp_args:Nc \__fp_tmp:w { \__fp_parse_infix_):N }

```

(End definition for __fp_parse_infix_):N.)

```

\__fp_parse_infix_,:N
\__fp_parse_infix_comma:w
\__fp_parse_apply_comma:NwNwN

```

As for other infix operations, if the previous operations has higher precedence the comma waits. Otherwise we call __fp_parse_operand:Nw to read more comma-delimited arguments that __fp_parse_infix_comma:w simply concatenates into a @-delimited array. The first comma in a tuple that is not a function argument is distinguished: in that case call __fp_parse_apply_comma:NwNwN whose job is to convert the first item of the tuple and an array of the remaining items into a tuple. In contrast to __fp_parse_apply_binary:NwNwN this function's operands are not single-object arrays.

```

18140 \cs_set_protected:Npn \__fp_tmp:w #1
18141 {
18142   \cs_new:Npn #1 ##1
18143   {
18144     \if_int_compare:w ##1 > \c__fp_prec_comma_int
18145       \exp_after:wN @
18146       \exp_after:wN \use_none:n
18147       \exp_after:wN #1
18148     \else:
18149       \if_int_compare:w ##1 < \c__fp_prec_comma_int
18150         \exp_after:wN @
18151         \exp_after:wN \__fp_parse_apply_comma:NwNwN
18152         \exp_after:wN ,
18153         \exp:w
18154       \else:
18155         \exp_after:wN \__fp_parse_infix_comma:w
18156         \exp:w
18157       \fi:
18158       \__fp_parse_operand:Nw \c__fp_prec_comma_int
18159       \exp_after:wN \__fp_parse_expand:w
18160     \fi:
18161   }
18162 }
18163 \exp_args:Nc \__fp_tmp:w { \__fp_parse_infix_,:N }
18164 \cs_new:Npn \__fp_parse_infix_comma:w #1 @
18165 { #1 @ \use_none:n }
18166 \cs_new:Npn \__fp_parse_apply_comma:NwNwN #1 #2@ #3 #4@ #5
18167 {
18168   \exp_after:wN \__fp_parse_continue:NwN

```

```

18169 \exp_after:wN #1
18170 \exp:w \exp_end_continue_f:w
18171 \__fp_exp_after_tuple_f:nw { }
18172 \s__fp_tuple \__fp_tuple_chk:w { #2 #4 } ;
18173 #5 #1
18174 }

```

(End definition for __fp_parse_infix_.:N, __fp_parse_infix_comma:w, and __fp_parse_apply_comma:NwNwN.)

28.7.2 Usual infix operators

As described in the “work plan”, each infix operator has an associated \..._infix... function, a computing function, and precedence, given as arguments to __fp_tmp:w. Using the general mechanism for arithmetic operations. The power operation must be associative in the opposite order from all others. For this, we use two distinct precedences.

```

\__fp_parse_infix_+:N
\__fp_parse_infix_-:N
\__fp_parse_infix_juxt:N
\__fp_parse_infix_/:N
\__fp_parse_infix_mul:N
\__fp_parse_infix_and:N
\__fp_parse_infix_or:N
\__fp_parse_infix_^:N
18175 \cs_set_protected:Npn \__fp_tmp:w #1#2#3#4
18176 {
18177   \cs_new:Npn #1 ##1
18178   {
18179     \if_int_compare:w ##1 < #3
18180       \exp_after:wN @
18181       \exp_after:wN \__fp_parse_apply_binary:NwNwN
18182       \exp_after:wN #2
18183       \exp:w
18184       \__fp_parse_operand:Nw #4
18185       \exp_after:wN \__fp_parse_expand:w
18186     \else:
18187       \exp_after:wN @
18188       \exp_after:wN \use_none:n
18189       \exp_after:wN #1
18190     \fi:
18191   }
18192 }
18193 \exp_args:Nc \__fp_tmp:w { \__fp_parse_infix_^:N } ^
18194 \c__fp_prec_hatii_int \c__fp_prec_hat_int
18195 \exp_args:Nc \__fp_tmp:w { \__fp_parse_infix_juxt:N } *
18196 \c__fp_prec_juxt_int \c__fp_prec_juxt_int
18197 \exp_args:Nc \__fp_tmp:w { \__fp_parse_infix_/:N } /
18198 \c__fp_prec_times_int \c__fp_prec_times_int
18199 \exp_args:Nc \__fp_tmp:w { \__fp_parse_infix_mul:N } *
18200 \c__fp_prec_times_int \c__fp_prec_times_int
18201 \exp_args:Nc \__fp_tmp:w { \__fp_parse_infix_-:N } -
18202 \c__fp_prec_plus_int \c__fp_prec_plus_int
18203 \exp_args:Nc \__fp_tmp:w { \__fp_parse_infix_+:N } +
18204 \c__fp_prec_plus_int \c__fp_prec_plus_int
18205 \exp_args:Nc \__fp_tmp:w { \__fp_parse_infix_and:N } &
18206 \c__fp_prec_and_int \c__fp_prec_and_int
18207 \exp_args:Nc \__fp_tmp:w { \__fp_parse_infix_or:N } |
18208 \c__fp_prec_or_int \c__fp_prec_or_int

```

(End definition for __fp_parse_infix_+:N and others.)

28.7.3 Juxtaposition

`__fp_parse_infix_(:N` When an opening parenthesis appears where we expect an infix operator, we compute the product of the previous operand and the contents of the parentheses using `__fp_parse_infix_mul:N`.

```
18209 \cs_new:cpn { __fp_parse_infix_(:N } #1
18210 { \__fp_parse_infix_mul:N #1 ( }
```

(End definition for `__fp_parse_infix_(:N`.)

28.7.4 Multi-character cases

`__fp_parse_infix_*:N`

```
18211 \cs_set_protected:Npn \__fp_tmp:w #1
18212 {
18213   \cs_new:cpn { __fp_parse_infix_*:N } ##1##2
18214   {
18215     \if:w * \exp_not:N ##2
18216       \exp_after:wN #1
18217       \exp_after:wN ##1
18218     \else:
18219       \exp_after:wN \__fp_parse_infix_mul:N
18220       \exp_after:wN ##1
18221       \exp_after:wN ##2
18222     \fi:
18223   }
18224 }
18225 \exp_args:Nc \__fp_tmp:w { __fp_parse_infix^:N }
```

(End definition for `__fp_parse_infix_*:N`.)

`__fp_parse_infix_|:Nw`

`__fp_parse_infix_&:Nw`

```
18226 \cs_set_protected:Npn \__fp_tmp:w #1#2#3
18227 {
18228   \cs_new:Npn #1 ##1##2
18229   {
18230     \if:w #2 \exp_not:N ##2
18231       \exp_after:wN #1
18232       \exp_after:wN ##1
18233       \exp:w \exp_after:wN \__fp_parse_expand:w
18234     \else:
18235       \exp_after:wN #3
18236       \exp_after:wN ##1
18237       \exp_after:wN ##2
18238     \fi:
18239   }
18240 }
18241 \exp_args:Nc \__fp_tmp:w { __fp_parse_infix_|:N } | \__fp_parse_infix_or:N
18242 \exp_args:Nc \__fp_tmp:w { __fp_parse_infix_&:N } & \__fp_parse_infix_and:N
```

(End definition for `__fp_parse_infix_|:Nw` and `__fp_parse_infix_&:Nw`.)

28.7.5 Ternary operator

```

__fp_parse_infix_?:N
__fp_parse_infix_:N
18243 \cs_set_protected:Npn \__fp_tmp:w #1#2#3#4
18244 {
18245   \cs_new:Npn #1 ##1
18246   {
18247     \if_int_compare:w ##1 < \c__fp_prec_quest_int
18248     #4
18249     \exp_after:wN @
18250     \exp_after:wN #2
18251     \exp:w
18252     \__fp_parse_operand:Nw #3
18253     \exp_after:wN \__fp_parse_expand:w
18254   \else:
18255     \exp_after:wN @
18256     \exp_after:wN \use_none:n
18257     \exp_after:wN #1
18258   \fi:
18259 }
18260 }
18261 \exp_args:Nc \__fp_tmp:w { __fp_parse_infix_?:N }
18262 \__fp_ternary:NwwN \c__fp_prec_quest_int { }
18263 \exp_args:Nc \__fp_tmp:w { __fp_parse_infix_:N }
18264 \__fp_ternary_auxii:NwwN \c__fp_prec_colon_int
18265 {
18266   \__kernel_msg_expandable_error:nnnn
18267   { kernel } { fp-missing } { ? } { ~for~?: }
18268 }

```

(End definition for __fp_parse_infix_?:N and __fp_parse_infix_:N.)

28.7.6 Comparisons

```

__fp_parse_infix_<:N
__fp_parse_infix_=:N
__fp_parse_infix_>:N
__fp_parse_infix_!:N
__fp_parse_excl_error:
__fp_parse_compare:NNNNNNN
  __fp_parse_compare_auxi:NNNNNNN
  __fp_parse_compare_auxii:NNNNN
  __fp_parse_compare_end:NNNNw
  __fp_compare:wNNNNw
18269 \cs_new:cpn { __fp_parse_infix_<:N } #1
18270 { \__fp_parse_compare:NNNNNNN #1 1 0 0 0 0 < }
18271 \cs_new:cpn { __fp_parse_infix_=:N } #1
18272 { \__fp_parse_compare:NNNNNNN #1 1 0 0 0 0 = }
18273 \cs_new:cpn { __fp_parse_infix_>:N } #1
18274 { \__fp_parse_compare:NNNNNNN #1 1 0 0 0 0 > }
18275 \cs_new:cpn { __fp_parse_infix_!:N } #1
18276 {
18277   \exp_after:wN \__fp_parse_compare:NNNNNNN
18278   \exp_after:wN #1
18279   \exp_after:wN 0
18280   \exp_after:wN 1
18281   \exp_after:wN 1
18282   \exp_after:wN 1
18283   \exp_after:wN 1
18284 }
18285 \cs_new:Npn \__fp_parse_excl_error:
18286 {
18287   \__kernel_msg_expandable_error:nnnn

```

```

18288         { kernel } { fp-missing } { = } { ~after~!. }
18289     }
18290 \cs_new:Npn \__fp_parse_compare:NNNNNNN #1
18291 {
18292     \if_int_compare:w #1 < \c__fp_prec_comp_int
18293         \exp_after:wN \__fp_parse_compare_auxi:NNNNNNN
18294         \exp_after:wN \__fp_parse_excl_error:
18295     \else:
18296         \exp_after:wN @
18297         \exp_after:wN \use_none:n
18298         \exp_after:wN \__fp_parse_compare:NNNNNNN
18299     \fi:
18300 }
18301 \cs_new:Npn \__fp_parse_compare_auxi:NNNNNNN #1#2#3#4#5#6#7
18302 {
18303     \if_case:w
18304         \__fp_int_eval:w \exp_after:wN ' \token_to_str:N #7 - '<
18305         \__fp_int_eval_end:
18306         \__fp_parse_compare_auxii:NNNNN #2#2#4#5#6
18307     \or: \__fp_parse_compare_auxii:NNNNN #2#3#2#5#6
18308     \or: \__fp_parse_compare_auxii:NNNNN #2#3#4#2#6
18309     \or: \__fp_parse_compare_auxii:NNNNN #2#3#4#5#2
18310     \else: #1 \__fp_parse_compare_end:NNNNw #3#4#5#6#7
18311     \fi:
18312 }
18313 \cs_new:Npn \__fp_parse_compare_auxii:NNNNN #1#2#3#4#5
18314 {
18315     \exp_after:wN \__fp_parse_compare_auxi:NNNNNNN
18316     \exp_after:wN \prg_do_nothing:
18317     \exp_after:wN #1
18318     \exp_after:wN #2
18319     \exp_after:wN #3
18320     \exp_after:wN #4
18321     \exp_after:wN #5
18322     \exp:w \exp_after:wN \__fp_parse_expand:w
18323 }
18324 \cs_new:Npn \__fp_parse_compare_end:NNNNw #1#2#3#4#5 \fi:
18325 {
18326     \fi:
18327     \exp_after:wN @
18328     \exp_after:wN \__fp_parse_apply_compare:NwNNNNNNwN
18329     \exp_after:wN \c_one_fp
18330     \exp_after:wN #1
18331     \exp_after:wN #2
18332     \exp_after:wN #3
18333     \exp_after:wN #4
18334     \exp:w
18335     \__fp_parse_operand:Nw \c__fp_prec_comp_int \__fp_parse_expand:w #5
18336 }
18337 \cs_new:Npn \__fp_parse_apply_compare:NwNNNNNNwN
18338     #1 #2@ #3 #4#5#6#7 #8@ #9
18339 {
18340     \if_int_odd:w
18341         \if_meaning:w \c_zero_fp #3

```

```

18342         0
18343     \else:
18344         \if_case:w \__fp_compare_back_any:ww #8 #2 \exp_stop_f:
18345             #5 \or: #6 \or: #7 \else: #4
18346         \fi:
18347     \fi:
18348     \exp_stop_f:
18349     \exp_after:wN \__fp_parse_apply_compare_aux:NNwN
18350     \exp_after:wN \c_one_fp
18351 \else:
18352     \exp_after:wN \__fp_parse_apply_compare_aux:NNwN
18353     \exp_after:wN \c_zero_fp
18354 \fi:
18355 #1 #8 #9
18356 }
18357 \cs_new:Npn \__fp_parse_apply_compare_aux:NNwN #1 #2 #3; #4
18358 {
18359     \if_meaning:w \__fp_parse_compare:NNNNNNN #4
18360         \exp_after:wN \__fp_parse_continue_compare:NNwNN
18361         \exp_after:wN #1
18362         \exp_after:wN #2
18363         \exp:w \exp_end_continue_f:w
18364         \__fp_exp_after_o:w #3;
18365         \exp:w \exp_end_continue_f:w
18366     \else:
18367         \exp_after:wN \__fp_parse_continue:NwN
18368         \exp_after:wN #2
18369         \exp:w \exp_end_continue_f:w
18370         \exp_after:wN #1
18371         \exp:w \exp_end_continue_f:w
18372     \fi:
18373     #4 #2
18374 }
18375 \cs_new:Npn \__fp_parse_continue_compare:NNwNN #1#2 #3@ #4#5
18376 { #4 #2 #3@ #1 }

```

(End definition for __fp_parse_infix_<:N and others.)

28.8 Tools for functions

__fp_parse_function_all_fp_o:fnw Followed by {<function name>} {<code>} <float array> @ this checks all floats are floating point numbers (no tuples).

```

18377 \cs_new:Npn \__fp_parse_function_all_fp_o:fnw #1#2#3 @
18378 {
18379     \__fp_array_if_all_fp:nTF {#3}
18380     { #2 #3 @ }
18381     {
18382         \__fp_error:nffn { fp-bad-args }
18383         {#1}
18384         { \fp_to_tl:n { \s__fp_tuple \__fp_tuple_chk:w {#3} ; } }
18385         { }
18386         \exp_after:wN \c_nan_fp
18387     }
18388 }

```

(End definition for `_fp_parse_function_all_fp_o:fnw`.)

`_fp_parse_function_one_two:nnw`
`_fp_parse_function_one_two_error_o:w`
`_fp_parse_function_one_two_aux:nnw`
`_fp_parse_function_one_two_auxii:nnw`

This is followed by `{(function name)}{(code)}{float array} @`. It checks that the `{float array}` consists of one or two floating point numbers (not tuples), then leaves the `{code}` (if there is one float) or its tail (if there are two floats) followed by the `{float array}`. The `{code}` should start with a single token such as `_fp_atan_default:w` that deals with the single-float case.

The first `_fp_if_type_fp:NTwFw` test catches the case of no argument and the case of a tuple argument. The next one distinguishes the case of a single argument (no error, just add `\c_one_fp`) from a tuple second argument. Finally check there is no further argument.

```

18389 \cs_new:Npn \_fp_parse_function_one_two:nnw #1#2#3
18390 {
18391   \_fp_if_type_fp:NTwFw
18392   #3 { } \s__fp \_fp_parse_function_one_two_error_o:w \s__fp_stop
18393   \_fp_parse_function_one_two_aux:nnw {#1} {#2} #3
18394 }
18395 \cs_new:Npn \_fp_parse_function_one_two_error_o:w #1#2#3#4 @
18396 {
18397   \_fp_error:nffn { fp-bad-args }
18398   {#2}
18399   { \fp_to_tl:n { \s__fp_tuple \_fp_tuple_chk:w {#4} ; } }
18400   { }
18401   \exp_after:wN \c_nan_fp
18402 }
18403 \cs_new:Npn \_fp_parse_function_one_two_aux:nnw #1#2 #3; #4
18404 {
18405   \_fp_if_type_fp:NTwFw
18406   #4 { }
18407   \s__fp
18408   {
18409     \if_meaning:w @ #4
18410     \exp_after:wN \use_iv:nnnn
18411     \fi:
18412     \_fp_parse_function_one_two_error_o:w
18413   }
18414   \s__fp_stop
18415   \_fp_parse_function_one_two_auxii:nnw {#1} {#2} #3; #4
18416 }
18417 \cs_new:Npn \_fp_parse_function_one_two_auxii:nnw #1#2#3; #4; #5
18418 {
18419   \if_meaning:w @ #5 \else:
18420     \exp_after:wN \_fp_parse_function_one_two_error_o:w
18421     \fi:
18422     \use_ii:nn {#1} { \use_none:n #2 } #3; #4; #5
18423 }

```

(End definition for `_fp_parse_function_one_two:nnw` and others.)

`_fp_tuple_map_o:nw`
`_fp_tuple_map_loop_o:nw`

Apply `#1` to all items in the following tuple and expand once afterwards. The code `#1` should itself expand once after its result.

```

18424 \cs_new:Npn \_fp_tuple_map_o:nw #1 \s__fp_tuple \_fp_tuple_chk:w #2 ;
18425 {

```



```

18426 \exp_after:wN \s__fp_tuple
18427 \exp_after:wN \__fp_tuple_chk:w
18428 \exp_after:wN {
18429 \exp:w \exp_end_continue_f:w
18430 \__fp_tuple_map_loop_o:nw {#1} #2
18431 { \s__fp \prg_break: } ;
18432 \prg_break_point:
18433 \exp_after:wN } \exp_after:wN ;
18434 }
18435 \cs_new:Npn \__fp_tuple_map_loop_o:nw #1#2#3 ;
18436 {
18437 \use_none:n #2
18438 #1 #2 #3 ;
18439 \exp:w \exp_end_continue_f:w
18440 \__fp_tuple_map_loop_o:nw {#1}
18441 }

```

(End definition for __fp_tuple_map_o:nw and __fp_tuple_map_loop_o:nw.)

__fp_tuple_mapthread_o:nww Apply #1 to pairs of items in the two following tuples and expand once afterwards.

```

\__fp_tuple_mapthread_loop_o:nw
18442 \cs_new:Npn \__fp_tuple_mapthread_o:nww #1
18443 \s__fp_tuple \__fp_tuple_chk:w #2 ;
18444 \s__fp_tuple \__fp_tuple_chk:w #3 ;
18445 {
18446 \exp_after:wN \s__fp_tuple
18447 \exp_after:wN \__fp_tuple_chk:w
18448 \exp_after:wN {
18449 \exp:w \exp_end_continue_f:w
18450 \__fp_tuple_mapthread_loop_o:nw {#1}
18451 #2 { \s__fp \prg_break: } ; @
18452 #3 { \s__fp \prg_break: } ;
18453 \prg_break_point:
18454 \exp_after:wN } \exp_after:wN ;
18455 }
18456 \cs_new:Npn \__fp_tuple_mapthread_loop_o:nw #1#2#3 ; #4 @ #5#6 ;
18457 {
18458 \use_none:n #2
18459 \use_none:n #5
18460 #1 #2 #3 ; #5 #6 ;
18461 \exp:w \exp_end_continue_f:w
18462 \__fp_tuple_mapthread_loop_o:nw {#1} #4 @
18463 }

```

(End definition for __fp_tuple_mapthread_o:nww and __fp_tuple_mapthread_loop_o:nw.)

28.9 Messages

```

18464 \__kernel_msg_new:nnn { kernel } { fp-deprecated }
18465 { '#1'~deprecated;~use~'#2' }
18466 \__kernel_msg_new:nnn { kernel } { unknown-fp-word }
18467 { Unknown~fp~word~#1. }
18468 \__kernel_msg_new:nnn { kernel } { fp-missing }
18469 { Missing~#1~inserted #2. }
18470 \__kernel_msg_new:nnn { kernel } { fp-extra }
18471 { Extra~#1~ignored. }

```

```

18472 \__kernel_msg_new:nnn { kernel } { fp-early-end }
18473 { Premature-end-in-fp-expression. }
18474 \__kernel_msg_new:nnn { kernel } { fp-after-e }
18475 { Cannot~use~#1 after~'e'. }
18476 \__kernel_msg_new:nnn { kernel } { fp-missing-number }
18477 { Missing-number~before~'#1'. }
18478 \__kernel_msg_new:nnn { kernel } { fp-unknown-symbol }
18479 { Unknown-symbol~#1~ignored. }
18480 \__kernel_msg_new:nnn { kernel } { fp-extra-comma }
18481 { Unexpected~comma~turned~to~nan~result. }
18482 \__kernel_msg_new:nnn { kernel } { fp-no-arg }
18483 { #1~got~no~argument;~used~nan. }
18484 \__kernel_msg_new:nnn { kernel } { fp-multi-arg }
18485 { #1~got~more~than~one~argument;~used~nan. }
18486 \__kernel_msg_new:nnn { kernel } { fp-num-args }
18487 { #1~expects~between~#2~and~#3~arguments. }
18488 \__kernel_msg_new:nnn { kernel } { fp-bad-args }
18489 { Arguments~in~#1#2~are~invalid. }
18490 \__kernel_msg_new:nnn { kernel } { fp-infty-pi }
18491 { Math-command~#1 is~not~an~fp }
18492 \cs_if_exist:cT { @unexpandable@protect }
18493 {
18494   \__kernel_msg_new:nnn { kernel } { fp-robust-cmd }
18495   { Robust~command~#1 invalid~in~fp-expression! }
18496 }
18497 </package>

```

29 l3fp-assign implementation

```

18498 <*package>
18499 <@@=fp>

```

29.1 Assigning values

\fp_new:N Floating point variables are initialized to be +0.

```

18500 \cs_new_protected:Npn \fp_new:N #1
18501 { \cs_new_eq:NN #1 \c_zero_fp }
18502 \cs_generate_variant:Nn \fp_new:N {c}

```

(End definition for \fp_new:N. This function is documented on page 204.)

\fp_set:Nn Simply use __fp_parse:n within various f-expanding assignments.

```

\fp_set:cn 18503 \cs_new_protected:Npn \fp_set:Nn #1#2
\fp_gset:Nn 18504 { \__kernel_tl_set:Nx #1 { \exp_not:f { \__fp_parse:n {#2} } } }
\fp_gset:cn 18505 \cs_new_protected:Npn \fp_gset:Nn #1#2
\fp_const:Nn 18506 { \__kernel_tl_gset:Nx #1 { \exp_not:f { \__fp_parse:n {#2} } } }
\fp_const:cn 18507 \cs_new_protected:Npn \fp_const:Nn #1#2
18508 { \tl_const:Nx #1 { \exp_not:f { \__fp_parse:n {#2} } } }
18509 \cs_generate_variant:Nn \fp_set:Nn {c}
18510 \cs_generate_variant:Nn \fp_gset:Nn {c}
18511 \cs_generate_variant:Nn \fp_const:Nn {c}

```

(End definition for \fp_set:Nn, \fp_gset:Nn, and \fp_const:Nn. These functions are documented on page 205.)

```

\fp_set_eq:NN Copying a floating point is the same as copying the underlying token list.
\fp_set_eq:cN 18512 \cs_new_eq:NN \fp_set_eq:NN \tl_set_eq:NN
\fp_set_eq:Nc 18513 \cs_new_eq:NN \fp_gset_eq:NN \tl_gset_eq:NN
\fp_set_eq:cc 18514 \cs_generate_variant:Nn \fp_set_eq:NN { c , Nc , cc }
\fp_gset_eq:NN 18515 \cs_generate_variant:Nn \fp_gset_eq:NN { c , Nc , cc }
\fp_gset_eq:cN
\fp_gset_eq:Nc (End definition for \fp_set_eq:NN and \fp_gset_eq:NN. These functions are documented on page 205.)
\fp_gset_eq:cc
\fp_zero:cN Setting a floating point to zero: copy \c_zero_fp.
\fp_zero:c 18516 \cs_new_protected:Npn \fp_zero:N #1 { \fp_set_eq:NN #1 \c_zero_fp }
\fp_gzero:N 18517 \cs_new_protected:Npn \fp_gzero:N #1 { \fp_gset_eq:NN #1 \c_zero_fp }
\fp_gzero:c 18518 \cs_generate_variant:Nn \fp_zero:N { c }
18519 \cs_generate_variant:Nn \fp_gzero:N { c }

(End definition for \fp_zero:N and \fp_gzero:N. These functions are documented on page 204.)

\fp_zero_new:N Set the floating point to zero, or define it if needed.
\fp_zero_new:c 18520 \cs_new_protected:Npn \fp_zero_new:N #1
\fp_gzero_new:N 18521 { \fp_if_exist:NTF #1 { \fp_zero:N #1 } { \fp_new:N #1 } }
\fp_gzero_new:c 18522 \cs_new_protected:Npn \fp_gzero_new:N #1
18523 { \fp_if_exist:NTF #1 { \fp_gzero:N #1 } { \fp_new:N #1 } }
18524 \cs_generate_variant:Nn \fp_zero_new:N { c }
18525 \cs_generate_variant:Nn \fp_gzero_new:N { c }

(End definition for \fp_zero_new:N and \fp_gzero_new:N. These functions are documented on page 205.)

```

29.2 Updating values

These match the equivalent functions in l3int and l3skip.

```

\fp_add:Nn For the sake of error recovery we should not simply set #1 to #1 ± (#2): for instance, if #2
\fp_add:cn is 0)+2, the parsing error would be raised at the last closing parenthesis rather than at
\fp_gadd:Nn the closing parenthesis in the user argument. Thus we evaluate #2 instead of just putting
\fp_gadd:cn parentheses. As an optimization we use \__fp_parse:n rather than \fp_eval:n, which
\fp_sub:Nn would convert the result away from the internal representation and back.
\fp_sub:cn 18526 \cs_new_protected:Npn \fp_add:Nn { \__fp_add:NNNn \fp_set:Nn + }
\fp_gsub:Nn 18527 \cs_new_protected:Npn \fp_gadd:Nn { \__fp_add:NNNn \fp_gset:Nn + }
\fp_gsub:cn 18528 \cs_new_protected:Npn \fp_sub:Nn { \__fp_add:NNNn \fp_set:Nn - }
\__fp_add:NNNn 18529 \cs_new_protected:Npn \fp_gsub:Nn { \__fp_add:NNNn \fp_gset:Nn - }
18530 \cs_new_protected:Npn \__fp_add:NNNn #1#2#3#4
18531 { #1 #3 { #3 #2 \__fp_parse:n {#4} } }
18532 \cs_generate_variant:Nn \fp_add:Nn { c }
18533 \cs_generate_variant:Nn \fp_gadd:Nn { c }
18534 \cs_generate_variant:Nn \fp_sub:Nn { c }
18535 \cs_generate_variant:Nn \fp_gsub:Nn { c }

(End definition for \fp_add:Nn and others. These functions are documented on page 205.)

```

29.3 Showing values

`\fp_show:N` This shows the result of computing its argument by passing the right data to `\tl_show:n` or `\tl_log:n`.

`\fp_show:c`

`\fp_log:N` 18536 `\cs_new_protected:Npn \fp_show:N { __fp_show:NN \tl_show:n }`

`\fp_log:c` 18537 `\cs_generate_variant:Nn \fp_show:N { c }`

`__fp_show:NN` 18538 `\cs_new_protected:Npn \fp_log:N { __fp_show:NN \tl_log:n }`

18539 `\cs_generate_variant:Nn \fp_log:N { c }`

18540 `\cs_new_protected:Npn __fp_show:NN #1#2`

18541 `{`

18542 `__kernel_chk_defined:NT #2`

18543 `{ \exp_args:Nx #1 { \token_to_str:N #2 = \fp_to_tl:N #2 } }`

18544 `}`

(End definition for `\fp_show:N`, `\fp_log:N`, and `__fp_show:NN`. These functions are documented on page 212.)

`\fp_show:n` Use general tools.

`\fp_log:n` 18545 `\cs_new_protected:Npn \fp_show:n`

18546 `{ \msg_show_eval:Nn \fp_to_tl:n }`

18547 `\cs_new_protected:Npn \fp_log:n`

18548 `{ \msg_log_eval:Nn \fp_to_tl:n }`

(End definition for `\fp_show:n` and `\fp_log:n`. These functions are documented on page 212.)

29.4 Some useful constants and scratch variables

`\c_one_fp` Some constants.

`\c_e_fp` 18549 `\fp_const:Nn \c_e_fp { 2.718 2818 2845 9045 }`

18550 `\fp_const:Nn \c_one_fp { 1 }`

(End definition for `\c_one_fp` and `\c_e_fp`. These variables are documented on page 210.)

`\c_pi_fp` We simply round π to and $\pi/180$ to 16 significant digits.

`\c_one_degree_fp` 18551 `\fp_const:Nn \c_pi_fp { 3.141 5926 5358 9793 }`

18552 `\fp_const:Nn \c_one_degree_fp { 0.0 1745 3292 5199 4330 }`

(End definition for `\c_pi_fp` and `\c_one_degree_fp`. These variables are documented on page 211.)

`\l_tmpa_fp` Scratch variables are simply initialized there.

`\l_tmpb_fp` 18553 `\fp_new:N \l_tmpa_fp`

`\g_tmpa_fp` 18554 `\fp_new:N \l_tmpb_fp`

`\g_tmpb_fp` 18555 `\fp_new:N \g_tmpa_fp`

18556 `\fp_new:N \g_tmpb_fp`

(End definition for `\l_tmpa_fp` and others. These variables are documented on page 211.)

18557 `\</package>`

30 l3fp-logic Implementation

18558 $\langle *package \rangle$

18559 $\langle @@=fp \rangle$

$\backslash_fp_parse_word_max:N$
 $\backslash_fp_parse_word_min:N$

Those functions may receive a variable number of arguments.

18560 $\backslash cs_new:Npn \backslash_fp_parse_word_max:N$
 18561 $\{ \backslash_fp_parse_function:NNN \backslash_fp_minmax_o:Nw 2 \}$
 18562 $\backslash cs_new:Npn \backslash_fp_parse_word_min:N$
 18563 $\{ \backslash_fp_parse_function:NNN \backslash_fp_minmax_o:Nw 0 \}$

(End definition for $\backslash_fp_parse_word_max:N$ and $\backslash_fp_parse_word_min:N$.)

30.1 Syntax of internal functions

- $\backslash_fp_compare_npos:nwnw \{ \langle expo_1 \rangle \} \langle body_1 \rangle ; \{ \langle expo_2 \rangle \} \langle body_2 \rangle ;$
- $\backslash_fp_minmax_o:Nw \langle sign \rangle \langle floating\ point\ array \rangle$
- $\backslash_fp_not_o:w ? \langle floating\ point\ array \rangle$ (with one floating point number only)
- $\backslash_fp_ \& _o:ww \langle floating\ point \rangle \langle floating\ point \rangle$
- $\backslash_fp_ | _o:ww \langle floating\ point \rangle \langle floating\ point \rangle$
- $\backslash_fp_ternary:NwwN, \backslash_fp_ternary_auxi:NwwN, \backslash_fp_ternary_auxii:NwwN$ have to be understood.

30.2 Tests

$\backslash fp_if_exist_p:N$
 $\backslash fp_if_exist_p:c$
 $\backslash fp_if_exist:N\TF$
 $\backslash fp_if_exist:c\TF$

Copies of the cs functions defined in l3basics.

18564 $\backslash prg_new_eq_conditional:NNn \backslash fp_if_exist:N \backslash cs_if_exist:N \{ TF , T , F , p \}$
 18565 $\backslash prg_new_eq_conditional:NNn \backslash fp_if_exist:c \backslash cs_if_exist:c \{ TF , T , F , p \}$

(End definition for $\backslash fp_if_exist:N\TF$. This function is documented on page 207.)

$\backslash fp_if_nan_p:n$
 $\backslash fp_if_nan:n\TF$

Evaluate and check if the result is a floating point of the same kind as NaN.

18566 $\backslash prg_new_conditional:Npnn \backslash fp_if_nan:n \#1 \{ TF , T , F , p \}$
 18567 $\{$
 18568 $\quad \backslash if:w 3 \backslash exp_last_unbraced:Nf \backslash_fp_kind:w \{ \backslash_fp_parse:n \{ \#1 \} \}$
 18569 $\quad \backslash prg_return_true:$
 18570 $\quad \backslash else:$
 18571 $\quad \backslash prg_return_false:$
 18572 $\quad \backslash fi:$
 18573 $\}$

(End definition for $\backslash fp_if_nan:n\TF$. This function is documented on page 269.)

30.3 Comparison

`\fp_compare_p:n`

`\fp_compare:nTF`

Within floating point expressions, comparison operators are treated as operations, so we evaluate #1, then compare with ± 0 . Tuples are true.

`__fp_compare_return:w`

```

18574 \prg_new_conditional:Npnn \fp_compare:n #1 { p , T , F , TF }
18575 {
18576   \exp_after:wN \__fp_compare_return:w
18577   \exp:w \exp_end_continue_f:w \__fp_parse:n {#1}
18578 }
18579 \cs_new:Npn \__fp_compare_return:w #1#2#3;
18580 {
18581   \if_charcode:w 0
18582     \__fp_if_type_fp:NTwFw
18583     #1 { \__fp_use_i_delimit_by_s_stop:nw #3 \s__fp_stop }
18584     \s__fp 1 \s__fp_stop
18585     \prg_return_false:
18586   \else:
18587     \prg_return_true:
18588   \fi:
18589 }
```

(End definition for `\fp_compare:nTF` and `__fp_compare_return:w`. This function is documented on page 208.)

`\fp_compare_p:nNn`

`\fp_compare:nNnTF`

`__fp_compare_aux:wn`

Evaluate #1 and #3, using an auxiliary to expand both, and feed the two floating point numbers swapped to `__fp_compare_back_any:ww`, defined below. Compare the result with ‘#2-‘=, which is -1 for $<$, 0 for $=$, 1 for $>$ and 2 for $?$.

```

18590 \prg_new_conditional:Npnn \fp_compare:nNn #1#2#3 { p , T , F , TF }
18591 {
18592   \if_int_compare:w
18593     \exp_after:wN \__fp_compare_aux:wn
18594     \exp:w \exp_end_continue_f:w \__fp_parse:n {#1} {#3}
18595     = \__fp_int_eval:w ‘#2 - ‘= \__fp_int_eval_end:
18596     \prg_return_true:
18597   \else:
18598     \prg_return_false:
18599   \fi:
18600 }
18601 \cs_new:Npn \__fp_compare_aux:wn #1; #2
18602 {
18603   \exp_after:wN \__fp_compare_back_any:ww
18604   \exp:w \exp_end_continue_f:w \__fp_parse:n {#2} #1;
18605 }
```

(End definition for `\fp_compare:nNnTF` and `__fp_compare_aux:wn`. This function is documented on page 207.)

`__fp_compare_back_any:ww`

`__fp_compare_back:ww`

`__fp_compare_nan:w`

`__fp_compare_back_any:ww` $\langle y \rangle$; $\langle x \rangle$;

Expands (in the same way as `\int_eval:n`) to -1 if $x < y$, 0 if $x = y$, 1 if $x > y$, and 2 otherwise (denoted as $x?y$). If either operand is `nan`, stop the comparison with `__fp_compare_nan:w` returning 2 . If x is negative, swap the outputs 1 and -1 (i.e., $>$ and $<$); we can henceforth assume that $x \geq 0$. If $y \geq 0$, and they have the same type, either they are normal and we compare them with `__fp_compare_npos:wnnw`, or they

are equal. If $y \geq 0$, but of a different type, the highest type is a larger number. Finally, if $y \leq 0$, then $x > y$, unless both are zero.

```

18606 \cs_new:Npn \__fp_compare_back_any:ww #1#2; #3
18607 {
18608   \__fp_if_type_fp:NTwFw
18609   #1 { \__fp_if_type_fp:NTwFw #3 \use_i:nn \s__fp \use_ii:nn \s__fp_stop }
18610   \s__fp \use_ii:nn \s__fp_stop
18611   \__fp_compare_back:ww
18612   {
18613     \cs:w
18614     __fp
18615     \__fp_type_from_scan:N #1
18616     _compare_back
18617     \__fp_type_from_scan:N #3
18618     :ww
18619     \cs_end:
18620   }
18621   #1#2 ; #3
18622 }
18623 \cs_new:Npn \__fp_compare_back:ww
18624   \s__fp \__fp_chk:w #1 #2 #3;
18625   \s__fp \__fp_chk:w #4 #5 #6;
18626 {
18627   \int_value:w
18628   \if_meaning:w 3 #1 \exp_after:wN \__fp_compare_nan:w \fi:
18629   \if_meaning:w 3 #4 \exp_after:wN \__fp_compare_nan:w \fi:
18630   \if_meaning:w 2 #5 - \fi:
18631   \if_meaning:w #2 #5
18632   \if_meaning:w #1 #4
18633   \if_meaning:w 1 #1
18634   \__fp_compare_npos:nwnw #6; #3;
18635   \else:
18636     0
18637   \fi:
18638   \else:
18639     \if_int_compare:w #4 < #1 - \fi: 1
18640   \fi:
18641   \else:
18642     \if_int_compare:w #1#4 = 0 \exp_stop_f:
18643     0
18644   \else:
18645     1
18646   \fi:
18647   \fi:
18648   \exp_stop_f:
18649 }
18650 \cs_new:Npn \__fp_compare_nan:w #1 \fi: \exp_stop_f: { 2 \exp_stop_f: }

```

(End definition for __fp_compare_back_any:ww, __fp_compare_back:ww, and __fp_compare_nan:w.)

__fp_compare_back_tuple:ww Tuple and floating point numbers are not comparable so return 2 in mixed cases or
 __fp_tuple_compare_back:ww when tuples have a different number of items. Otherwise compare pairs of items with
 __fp_tuple_compare_back_tuple:ww __fp_compare_back_any:ww and if any don't match return 2 (as \int_value:w 02
 __fp_tuple_compare_back_loop:w \exp_stop_f:).

```

18651 \cs_new:Npn \__fp_compare_back_tuple:ww #1; #2; { 2 }
18652 \cs_new:Npn \__fp_tuple_compare_back:ww #1; #2; { 2 }
18653 \cs_new:Npn \__fp_tuple_compare_back_tuple:ww
18654   \s__fp_tuple \__fp_tuple_chk:w #1;
18655   \s__fp_tuple \__fp_tuple_chk:w #2;
18656   {
18657     \int_compare:nNnTF { \__fp_array_count:n {#1} } =
18658       { \__fp_array_count:n {#2} }
18659       {
18660         \int_value:w 0
18661         \__fp_tuple_compare_back_loop:w
18662           #1 { \s__fp \prg_break: } ; @
18663           #2 { \s__fp \prg_break: } ;
18664         \prg_break_point:
18665         \exp_stop_f:
18666       }
18667       { 2 }
18668   }
18669 \cs_new:Npn \__fp_tuple_compare_back_loop:w #1#2 ; #3 @ #4#5 ;
18670 {
18671   \use_none:n #1
18672   \use_none:n #4
18673   \if_int_compare:w
18674     \__fp_compare_back_any:ww #1 #2 ; #4 #5 ; = 0 \exp_stop_f:
18675   \else:
18676     2 \exp_after:wN \prg_break:
18677   \fi:
18678   \__fp_tuple_compare_back_loop:w #3 @
18679 }

```

(End definition for __fp_compare_back_tuple:ww and others.)

__fp_compare_npos:nwnw
 __fp_compare_significand:nnnnnnnn

__fp_compare_npos:nwnw {<expo₁>} <body₁> ; {<expo₂>} <body₂> ;
 Within an \int_value:w ... \exp_stop_f: construction, this expands to 0 if the two numbers are equal, -1 if the first is smaller, and 1 if the first is bigger. First compare the exponents: the larger one denotes the larger number. If they are equal, we must compare significands. If both the first 8 digits and the next 8 digits coincide, the numbers are equal. If only the first 8 digits coincide, the next 8 decide. Otherwise, the first 8 digits are compared.

```

18680 \cs_new:Npn \__fp_compare_npos:nwnw #1#2; #3#4;
18681 {
18682   \if_int_compare:w #1 = #3 \exp_stop_f:
18683     \__fp_compare_significand:nnnnnnnn #2 #4
18684   \else:
18685     \if_int_compare:w #1 < #3 - \fi: 1
18686   \fi:
18687 }
18688 \cs_new:Npn \__fp_compare_significand:nnnnnnnn #1#2#3#4#5#6#7#8
18689 {
18690   \if_int_compare:w #1#2 = #5#6 \exp_stop_f:
18691   \if_int_compare:w #3#4 = #7#8 \exp_stop_f:
18692     0
18693   \else:
18694     \if_int_compare:w #3#4 < #7#8 - \fi: 1

```



```

18695     \fi:
18696   \else:
18697     \if_int_compare:w #1#2 < #5#6 - \fi: 1
18698   \fi:
18699 }

```

(End definition for `__fp_compare_npos:nwnw` and `__fp_compare_significand:nnnnnnnn`.)

30.4 Floating point expression loops

`\fp_do_until:nn` These are quite easy given the above functions. The `do_until` and `do_while` versions execute the body, then test. The `until_do` and `while_do` do it the other way round.

```

18700 \cs_new:Npn \fp_do_until:nn #1#2
18701 {
18702   #2
18703   \fp_compare:nF {#1}
18704   { \fp_do_until:nn {#1} {#2} }
18705 }
18706 \cs_new:Npn \fp_do_while:nn #1#2
18707 {
18708   #2
18709   \fp_compare:nT {#1}
18710   { \fp_do_while:nn {#1} {#2} }
18711 }
18712 \cs_new:Npn \fp_until_do:nn #1#2
18713 {
18714   \fp_compare:nF {#1}
18715   {
18716     #2
18717     \fp_until_do:nn {#1} {#2}
18718   }
18719 }
18720 \cs_new:Npn \fp_while_do:nn #1#2
18721 {
18722   \fp_compare:nT {#1}
18723   {
18724     #2
18725     \fp_while_do:nn {#1} {#2}
18726   }
18727 }

```

(End definition for `\fp_do_until:nn` and others. These functions are documented on page 209.)

`\fp_do_until:nNnn` As above but not using the `nNn` syntax.

```

18728 \cs_new:Npn \fp_do_until:nNnn #1#2#3#4
18729 {
18730   #4
18731   \fp_compare:nNnF {#1} #2 {#3}
18732   { \fp_do_until:nNnn {#1} #2 {#3} {#4} }
18733 }
18734 \cs_new:Npn \fp_do_while:nNnn #1#2#3#4
18735 {
18736   #4
18737   \fp_compare:nNnT {#1} #2 {#3}

```

```

18738     { \fp_do_while:nNnn {#1} #2 {#3} {#4} }
18739   }
18740 \cs_new:Npn \fp_until_do:nNnn #1#2#3#4
18741 {
18742   \fp_compare:nNnF {#1} #2 {#3}
18743   {
18744     #4
18745     \fp_until_do:nNnn {#1} #2 {#3} {#4}
18746   }
18747 }
18748 \cs_new:Npn \fp_while_do:nNnn #1#2#3#4
18749 {
18750   \fp_compare:nNnT {#1} #2 {#3}
18751   {
18752     #4
18753     \fp_while_do:nNnn {#1} #2 {#3} {#4}
18754   }
18755 }

```

(End definition for `\fp_do_until:nNnn` and others. These functions are documented on page 208.)

\fp_step_function:nnnN

\fp_step_function:nnnc

`__fp_step:wwwN`

`__fp_step_fp:wwwN`

`__fp_step:NnnnnN`

`__fp_step:NfnnnN`

The approach here is somewhat similar to `\int_step_function:nnnN`. There are two subtleties: we use the internal parser `__fp_parse:n` to avoid converting back and forth from the internal representation; and (due to rounding) even a non-zero step does not guarantee that the loop counter increases.

```

18756 \cs_new:Npn \fp_step_function:nnnN #1#2#3
18757 {
18758   \exp_after:wN \__fp_step:wwwN
18759   \exp:w \exp_end_continue_f:w \__fp_parse_o:n {#1}
18760   \exp:w \exp_end_continue_f:w \__fp_parse_o:n {#2}
18761   \exp:w \exp_end_continue_f:w \__fp_parse:n {#3}
18762 }
18763 \cs_generate_variant:Nn \fp_step_function:nnnN { nnnc }
18764 % \end{macrocode}
18765 % Only floating point numbers (not tuples) are allowed arguments.
18766 % Only \enquote{normal} floating points (not $\pm 0$,
18767 % $\pm\texttt{inf}$, $\texttt{nan}$) can be used as step; if positive,
18768 % call \cs{__fp_step:NnnnnN} with argument |>| otherwise~|<|. This
18769 % function has one more argument than its integer counterpart, namely
18770 % the previous value, to catch the case where the loop has made no
18771 % progress. Conversion to decimal is done just before calling the
18772 % user's function.
18773 % \begin{macrocode}
18774 \cs_new:Npn \__fp_step:wwwN #1#2; #3#4; #5#6; #7
18775 {
18776   \__fp_if_type_fp:NTwFw #1 { } \s__fp \prg_break: \s__fp_stop
18777   \__fp_if_type_fp:NTwFw #3 { } \s__fp \prg_break: \s__fp_stop
18778   \__fp_if_type_fp:NTwFw #5 { } \s__fp \prg_break: \s__fp_stop
18779   \use_i:nnnn { \__fp_step_fp:wwwN #1#2; #3#4; #5#6; #7 }
18780   \prg_break_point:
18781   \use:n
18782   {
18783     \__fp_error:nfff { fp-step-tuple } { \fp_to_tl:n { #1#2 ; } }
18784     { \fp_to_tl:n { #3#4 ; } } { \fp_to_tl:n { #5#6 ; } }

```

```

18785     }
18786   }
18787   \cs_new:Npn \__fp_step_fp:wwwN #1 ; \s__fp \__fp_chk:w #2#3#4 ; #5; #6
18788   {
18789     \token_if_eq_meaning:NNTF #2 1
18790     {
18791       \token_if_eq_meaning:NNTF #3 0
18792       { \__fp_step:NnnnnN > }
18793       { \__fp_step:NnnnnN < }
18794     }
18795     {
18796       \token_if_eq_meaning:NNTF #2 0
18797       {
18798         \__kernel_msg_expandable_error:nnn { kernel }
18799         { zero-step } {#6}
18800       }
18801       {
18802         \__fp_error:nnfn { fp-bad-step } { }
18803         { \fp_to_tl:n { \s__fp \__fp_chk:w #2#3#4 ; } } {#6}
18804       }
18805       \use_none:nnnnn
18806     }
18807     { #1 ; } { \c_nan_fp } { \s__fp \__fp_chk:w #2#3#4 ; } { #5 ; } #6
18808   }
18809   \cs_new:Npn \__fp_step:NnnnnN #1#2#3#4#5#6
18810   {
18811     \fp_compare:nNnTF {#2} = {#3}
18812     {
18813       \__fp_error:nffn { fp-tiny-step }
18814       { \fp_to_tl:n {#3} } { \fp_to_tl:n {#4} } {#6}
18815     }
18816     {
18817       \fp_compare:nNnF {#2} #1 {#5}
18818       {
18819         \exp_args:Nf #6 { \__fp_to_decimal_dispatch:w #2 }
18820         \__fp_step:NfnnnnN
18821         #1 { \__fp_parse:n { #2 + #4 } } {#2} {#4} {#5} #6
18822       }
18823     }
18824   }
18825   \cs_generate_variant:Nn \__fp_step:NnnnnN { Nf }

```

(End definition for \fp_step_function:nnnN and others. This function is documented on page 210.)

\fp_step_inline:nnnn As for \int_step_inline:nnnn, create a global function and apply it, following up with
\fp_step_variable:nnnNn a break point.

```

\__fp_step:NNnnnn
18826 \cs_new_protected:Npn \fp_step_inline:nnnn
18827 {
18828   \int_gincr:N \g__kernel_prg_map_int
18829   \exp_args:NNc \__fp_step:NNnnnn
18830   \cs_gset_protected:Npn
18831   { __fp_map_ \int_use:N \g__kernel_prg_map_int :w }
18832 }
18833 \cs_new_protected:Npn \fp_step_variable:nnnNn #1#2#3#4#5

```

```

18834 {
18835   \int_gincr:N \g__kernel_prg_map_int
18836   \exp_args:Nnc \__fp_step:NNnnnn
18837   \cs_gset_protected:Npx
18838     { \__fp_map_ \int_use:N \g__kernel_prg_map_int :w }
18839     {#1} {#2} {#3}
18840     {
18841       \tl_set:Nn \exp_not:N #4 {##1}
18842       \exp_not:n {#5}
18843     }
18844   }
18845   \cs_new_protected:Npn \__fp_step:NNnnnn #1#2#3#4#5#6
18846   {
18847     #1 #2 ##1 {#6}
18848     \fp_step_function:nnnN {#3} {#4} {#5} #2
18849     \prg_break_point:Nn \scan_stop: { \int_gdecr:N \g__kernel_prg_map_int }
18850   }

```

(End definition for `\fp_step_inline:nnnn`, `\fp_step_variable:nnnN`, and `__fp_step:NNnnnn`. These functions are documented on page 210.)

```

18851 \__kernel_msg_new:nnn { kernel } { fp-step-tuple }
18852 { Tuple~argument~in~fp_step...~{#1}{#2}{#3}. }
18853 \__kernel_msg_new:nnn { kernel } { fp-bad-step }
18854 { Invalid~step~size~#2~in~step~function~#3. }
18855 \__kernel_msg_new:nnn { kernel } { fp-tiny-step }
18856 { Tiny~step~size~(#{1}+#{2}=#{1})~in~step~function~#3. }

```

30.5 Extrema

```

\__fp_minmax_o:Nw
\__fp_minmax_aux_o:Nw

```

First check all operands are floating point numbers. The argument #1 is 2 to find the maximum of an array #2 of floating point numbers, and 0 to find the minimum. We read numbers sequentially, keeping track of the largest (smallest) number found so far. If numbers are equal (for instance ± 0), the first is kept. We append $-\infty$ (∞), for the case of an empty array. Since no number is smaller (larger) than that, this additional item only affects the maximum (minimum) in the case of `max()` and `min()` with no argument. The weird fp-like trailing marker breaks the loop correctly: see the precise definition of `__fp_minmax_loop:Nww`.

```

18857 \cs_new:Npn \__fp_minmax_o:Nw #1
18858 {
18859   \__fp_parse_function_all_fp_o:fnw
18860   { \token_if_eq_meaning:NNTF 0 #1 { min } { max } }
18861   { \__fp_minmax_aux_o:Nw #1 }
18862 }
18863 \cs_new:Npn \__fp_minmax_aux_o:Nw #1#2 @
18864 {
18865   \if_meaning:w 0 #1
18866     \exp_after:wN \__fp_minmax_loop:Nww \exp_after:wN +
18867   \else:
18868     \exp_after:wN \__fp_minmax_loop:Nww \exp_after:wN -
18869   \fi:
18870   #2
18871   \s__fp \__fp_chk:w 2 #1 \s__fp_exact ;
18872   \s__fp \__fp_chk:w { 3 \__fp_minmax_break_o:w } ;

```

```
18873 }
```

(End definition for `_fp_minmax_o:Nw` and `_fp_minmax_aux_o:Nw`.)

`_fp_minmax_loop:Nww`

The first argument is `-` or `+` to denote the case where the currently largest (smallest) number found (first floating point argument) should be replaced by the new number (second floating point argument). If the new number is `nan`, keep that as the extremum, unless that extremum is already a `nan`. Otherwise, compare the two numbers. If the new number is larger (in the case of `max`) or smaller (in the case of `min`), the test yields `true`, and we keep the second number as a new maximum; otherwise we keep the first number. Then loop.

```
18874 \cs_new:Npn \_fp_minmax_loop:Nww
18875   #1 \s__fp \_fp_chk:w #2#3; \s__fp \_fp_chk:w #4#5;
18876   {
18877     \if_meaning:w 3 #4
18878     \if_meaning:w 3 #2
18879       \_fp_minmax_auxi:ww
18880     \else:
18881       \_fp_minmax_auxii:ww
18882     \fi:
18883   \else:
18884     \if_int_compare:w
18885       \_fp_compare_back:ww
18886       \s__fp \_fp_chk:w #4#5;
18887       \s__fp \_fp_chk:w #2#3;
18888       = #1 1 \exp_stop_f:
18889       \_fp_minmax_auxii:ww
18890     \else:
18891       \_fp_minmax_auxi:ww
18892     \fi:
18893   \fi:
18894   \_fp_minmax_loop:Nww #1
18895   \s__fp \_fp_chk:w #2#3;
18896   \s__fp \_fp_chk:w #4#5;
18897 }
```

(End definition for `_fp_minmax_loop:Nww`.)

`_fp_minmax_auxi:ww`

Keep the first/second number, and remove the other.

`_fp_minmax_auxii:ww`

```
18898 \cs_new:Npn \_fp_minmax_auxi:ww #1 \fi: \fi: #2 \s__fp #3 ; \s__fp #4;
18899   { \fi: \fi: #2 \s__fp #3 ; }
18900 \cs_new:Npn \_fp_minmax_auxii:ww #1 \fi: \fi: #2 \s__fp #3 ;
18901   { \fi: \fi: #2 }
```

(End definition for `_fp_minmax_auxi:ww` and `_fp_minmax_auxii:ww`.)

`_fp_minmax_break_o:w`

This function is called from within an `\if_meaning:w` test. Skip to the end of the tests, close the current test with `\fi:`, clean up, and return the appropriate number with one post-expansion.

```
18902 \cs_new:Npn \_fp_minmax_break_o:w #1 \fi: \fi: #2 \s__fp #3; #4;
18903   { \fi: \_fp_exp_after_o:w \s__fp #3; }
```

(End definition for `_fp_minmax_break_o:w`.)

30.6 Boolean operations

`__fp_not_o:w` Return true or false, with two expansions, one to exit the conditional, and one to please l3fp-parse. The first argument is provided by l3fp-parse and is ignored.

```

18904 \cs_new:Npn \__fp_not_o:w #1 \s__fp \__fp_chk:w #2#3; @
18905 {
18906   \if_meaning:w 0 #2
18907     \exp_after:wN \exp_after:wN \exp_after:wN \c_one_fp
18908   \else:
18909     \exp_after:wN \exp_after:wN \exp_after:wN \c_zero_fp
18910   \fi:
18911 }
18912 \cs_new:Npn \__fp_tuple_not_o:w #1 @ { \exp_after:wN \c_zero_fp }
```

(End definition for `__fp_not_o:w` and `__fp_tuple_not_o:w`.)

`__fp_&_o:ww` For and, if the first number is zero, return it (with the same sign). Otherwise, return the second one. For or, the logic is reversed: if the first number is non-zero, return it, otherwise return the second number: we achieve that by hi-jacking `__fp_&_o:ww`, inserting an extra argument, `\else:`, before `\s__fp`. In all cases, expand after the floating point number.

```

18913 \group_begin:
18914   \char_set_catcode_letter:N &
18915   \char_set_catcode_letter:N |
18916   \cs_new:Npn \__fp_&_o:ww #1 \s__fp \__fp_chk:w #2#3;
18917   {
18918     \if_meaning:w 0 #2 #1
18919       \__fp_and_return:wNw \s__fp \__fp_chk:w #2#3;
18920     \fi:
18921     \__fp_exp_after_o:w
18922   }
18923   \cs_new:Npn \__fp_&_tuple_o:ww #1 \s__fp \__fp_chk:w #2#3;
18924   {
18925     \if_meaning:w 0 #2 #1
18926       \__fp_and_return:wNw \s__fp \__fp_chk:w #2#3;
18927     \fi:
18928     \__fp_exp_after_tuple_o:w
18929   }
18930   \cs_new:Npn \__fp_tuple_&_o:ww #1; { \__fp_exp_after_o:w }
18931   \cs_new:Npn \__fp_tuple_&_tuple_o:ww #1; { \__fp_exp_after_tuple_o:w }
18932   \cs_new:Npn \__fp_|_o:ww { \__fp_&_o:ww \else: }
18933   \cs_new:Npn \__fp_|_tuple_o:ww { \__fp_&_tuple_o:ww \else: }
18934   \cs_new:Npn \__fp_tuple_|_o:ww #1; #2; { \__fp_exp_after_tuple_o:w #1; }
18935   \cs_new:Npn \__fp_tuple_|_tuple_o:ww #1; #2;
18936     { \__fp_exp_after_tuple_o:w #1; }
18937 \group_end:
18938 \cs_new:Npn \__fp_and_return:wNw #1; \fi: #2;
18939 { \fi: \__fp_exp_after_o:w #1; }
```

(End definition for `__fp_&_o:ww` and others.)

30.7 Ternary operator

`_fp_ternary:NwN`
`_fp_ternary_auxi:NwN`
`_fp_ternary_auxii:NwN`

The first function receives the test and the true branch of the `?:` ternary operator. It calls `_fp_ternary_auxii:NwN` if the test branch is a floating point number ± 0 , and otherwise calls `_fp_ternary_auxi:NwN`. These functions select one of their two arguments.

```

18940 \cs_new:Npn \_fp_ternary:NwN #1 #2#3@ #4@ #5
18941 {
18942   \if_meaning:w \_fp_parse_infix_:N #5
18943     \if_charcode:w 0
18944       \_fp_if_type_fp:NTwFw
18945       #2 { \use_i:nn \_fp_use_i_delimit_by_s_stop:nw #3 \s__fp_stop }
18946       \s__fp 1 \s__fp_stop
18947       \exp_after:wN \exp_after:wN \exp_after:wN \_fp_ternary_auxii:NwN
18948     \else:
18949       \exp_after:wN \exp_after:wN \exp_after:wN \_fp_ternary_auxi:NwN
18950     \fi:
18951     \exp_after:wN #1
18952     \exp:w \exp_end_continue_f:w
18953     \_fp_exp_after_array_f:w #4 \s__fp_expr_stop
18954     \exp_after:wN @
18955     \exp:w
18956     \_fp_parse_operand:Nw \c__fp_prec_colon_int
18957     \_fp_parse_expand:w
18958   \else:
18959     \__kernel_msg_expandable_error:nnnn
18960     { kernel } { fp-missing } { : } { ~for~?: }
18961     \exp_after:wN \_fp_parse_continue:NwN
18962     \exp_after:wN #1
18963     \exp:w \exp_end_continue_f:w
18964     \_fp_exp_after_array_f:w #4 \s__fp_expr_stop
18965     \exp_after:wN #5
18966     \exp_after:wN #1
18967   \fi:
18968 }
18969 \cs_new:Npn \_fp_ternary_auxi:NwN #1#2@#3@#4
18970 {
18971   \exp_after:wN \_fp_parse_continue:NwN
18972   \exp_after:wN #1
18973   \exp:w \exp_end_continue_f:w
18974   \_fp_exp_after_array_f:w #2 \s__fp_expr_stop
18975   #4 #1
18976 }
18977 \cs_new:Npn \_fp_ternary_auxii:NwN #1#2@#3@#4
18978 {
18979   \exp_after:wN \_fp_parse_continue:NwN
18980   \exp_after:wN #1
18981   \exp:w \exp_end_continue_f:w
18982   \_fp_exp_after_array_f:w #3 \s__fp_expr_stop
18983   #4 #1
18984 }

```

(End definition for `_fp_ternary:NwN`, `_fp_ternary_auxi:NwN`, and `_fp_ternary_auxii:NwN`.)

```

18985 \</package>

```

31 l3fp-basics Implementation

```
18986 ⟨*package⟩
18987 ⟨@@=fp⟩
```

The `l3fp-basics` module implements addition, subtraction, multiplication, and division of two floating points, and the absolute value and sign-changing operations on one floating point. All operations implemented in this module yield the outcome of rounding the infinitely precise result of the operation to the nearest floating point.

Some algorithms used below end up being quite similar to some described in “What Every Computer Scientist Should Know About Floating Point Arithmetic”, by David Goldberg, which can be found at <http://cr.yp.to/2005-590/goldberg.pdf>.

Unary functions.

```

__fp_parse_word_abs:N
__fp_parse_word_logb:N
__fp_parse_word_sign:N
__fp_parse_word_sqrt:N
18988 \cs_new:Npn \__fp_parse_word_abs:N
18989   { \__fp_parse_unary_function:NNN \__fp_set_sign_o:w 0 }
18990 \cs_new:Npn \__fp_parse_word_logb:N
18991   { \__fp_parse_unary_function:NNN \__fp_logb_o:w ? }
18992 \cs_new:Npn \__fp_parse_word_sign:N
18993   { \__fp_parse_unary_function:NNN \__fp_sign_o:w ? }
18994 \cs_new:Npn \__fp_parse_word_sqrt:N
18995   { \__fp_parse_unary_function:NNN \__fp_sqrt_o:w ? }
```

(End definition for `__fp_parse_word_abs:N` and others.)

31.1 Addition and subtraction

We define here two functions, `__fp_-_o:ww` and `__fp+_o:ww`, which perform the subtraction and addition of their two floating point operands, and expand the tokens following the result once.

A more obscure function, `__fp_add_big_i_o:wNww`, is used in `l3fp-expo`.

The logic goes as follows:

- `__fp_-_o:ww` calls `__fp+_o:ww` to do the work, with the sign of the second operand flipped;
- `__fp+_o:ww` dispatches depending on the type of floating point, calling specialized auxiliaries;
- in all cases except summing two normal floating point numbers, we return one or the other operands depending on the signs, or detect an invalid operation in the case of $\infty - \infty$;
- for normal floating point numbers, compare the signs;
- to add two floating point numbers of the same sign or of opposite signs, shift the significand of the smaller one to match the bigger one, perform the addition or subtraction of significands, check for a carry, round, and pack using the `__fp-basics_pack_...` functions.

The trickiest part is to round correctly when adding or subtracting normal floating point numbers.

31.1.1 Sign, exponent, and special numbers

`__fp_-_o:ww` The `__fp+_o:ww` auxiliary has a hook: it takes one argument between the first `\s__fp` and `__fp_chk:w`, which is applied to the sign of the second operand. Positioning the hook there means that `__fp+_o:ww` can still perform the sanity check that it was followed by `\s__fp`.

```
18996 \cs_new:cpx { __fp_-_o:ww } \s__fp
18997 {
18998   \exp_not:c { __fp+_o:ww }
18999   \exp_not:n { \s__fp \__fp_neg_sign:N }
19000 }
```

(End definition for `__fp_-_o:ww`.)

`__fp+_o:ww` This function is either called directly with an empty #1 to compute an addition, or it is called by `__fp_-_o:ww` with `__fp_neg_sign:N` as #1 to compute a subtraction, in which case the second operand's sign should be changed. If the *<types>* #2 and #4 are the same, dispatch to case #2 (0, 1, 2, or 3), where we call specialized functions: thanks to `\int_value:w`, those receive the tweaked *<sign₂>* (expansion of #1#5) as an argument. If the *<types>* are distinct, the result is simply the floating point number with the highest *<type>*. Since case 3 (used for two nan) also picks the first operand, we can also use it when *<type₁>* is greater than *<type₂>*. Also note that we don't need to worry about *<sign₂>* in that case since the second operand is discarded.

```
19001 \cs_new:cpn { __fp+_o:ww }
19002   \s__fp #1 \__fp_chk:w #2 #3 ; \s__fp \__fp_chk:w #4 #5
19003 {
19004   \if_case:w
19005     \if_meaning:w #2 #4
19006       #2
19007     \else:
19008       \if_int_compare:w #2 > #4 \exp_stop_f:
19009         3
19010       \else:
19011         4
19012       \fi:
19013     \fi:
19014     \exp_stop_f:
19015     \exp_after:wN \__fp_add_zeros_o:Nww \int_value:w
19016   \or:   \exp_after:wN \__fp_add_normal_o:Nww \int_value:w
19017   \or:   \exp_after:wN \__fp_add_inf_o:Nww \int_value:w
19018   \or:   \__fp_case_return_i_o:ww
19019   \else: \exp_after:wN \__fp_add_return_ii_o:Nww \int_value:w
19020   \fi:
19021   #1 #5
19022   \s__fp \__fp_chk:w #2 #3 ;
19023   \s__fp \__fp_chk:w #4 #5
19024 }
```

(End definition for `__fp+_o:ww`.)

`__fp_add_return_ii_o:Nww` Ignore the first operand, and return the second, but using the sign #1 rather than #4. As usual, expand after the floating point.

```
19025 \cs_new:Npn \__fp_add_return_ii_o:Nww #1 #2 ; \s__fp \__fp_chk:w #3 #4
19026 { \__fp_exp_after_o:w \s__fp \__fp_chk:w #3 #1 }
```

(End definition for _fp_add_return_ii_o:Nww.)

_fp_add_zeros_o:Nww Adding two zeros yields \c_zero_fp, except if both zeros were -0 .

```

19027 \cs_new:Npn \_fp_add_zeros_o:Nww #1 \s__fp \_fp_chk:w 0 #2
19028 {
19029   \if_int_compare:w #2 #1 = 20 \exp_stop_f:
19030     \exp_after:wN \_fp_add_return_ii_o:Nww
19031   \else:
19032     \_fp_case_return_i_o:ww
19033   \fi:
19034   #1
19035   \s__fp \_fp_chk:w 0 #2
19036 }

```

(End definition for _fp_add_zeros_o:Nww.)

_fp_add_inf_o:Nww If both infinities have the same sign, just return that infinity, otherwise, it is an invalid operation. We find out if that invalid operation is an addition or a subtraction by testing whether the tweaked $\langle sign_2 \rangle$ (#1) and the $\langle sign_2 \rangle$ (#4) are identical.

```

19037 \cs_new:Npn \_fp_add_inf_o:Nww
19038   #1 \s__fp \_fp_chk:w 2 #2 #3; \s__fp \_fp_chk:w 2 #4
19039 {
19040   \if_meaning:w #1 #2
19041     \_fp_case_return_i_o:ww
19042   \else:
19043     \_fp_case_use:nw
19044     {
19045       \exp_last_unbraced:Nf \_fp_invalid_operation_o:Nww
19046       { \token_if_eq_meaning:NNTF #1 #4 + - }
19047     }
19048   \fi:
19049   \s__fp \_fp_chk:w 2 #2 #3;
19050   \s__fp \_fp_chk:w 2 #4
19051 }

```

(End definition for _fp_add_inf_o:Nww.)

_fp_add_normal_o:Nww _fp_add_normal_o:Nww $\langle sign_2 \rangle$ \s__fp _fp_chk:w 1 $\langle sign_1 \rangle$ $\langle exp_1 \rangle$ $\langle body_1 \rangle$; \s__fp _fp_chk:w 1 $\langle initial\ sign_2 \rangle$ $\langle exp_2 \rangle$ $\langle body_2 \rangle$;

We now have two normal numbers to add, and we have to check signs and exponents more carefully before performing the addition.

```

19052 \cs_new:Npn \_fp_add_normal_o:Nww #1 \s__fp \_fp_chk:w 1 #2
19053 {
19054   \if_meaning:w #1#2
19055     \exp_after:wN \_fp_add_npos_o:NnwNnw
19056   \else:
19057     \exp_after:wN \_fp_sub_npos_o:NnwNnw
19058   \fi:
19059   #2
19060 }

```

(End definition for _fp_add_normal_o:Nww.)

31.1.2 Absolute addition

In this subsection, we perform the addition of two positive normal numbers.

$$\backslash_fp_add_npos_o:NnwNnw \langle sign_1 \rangle \langle exp_1 \rangle \langle body_1 \rangle ; \backslash s_fp \backslash_fp_chk:w 1$$

Since we are doing an addition, the final sign is $\langle sign_1 \rangle$. Start an `_fp_int_eval:w`, responsible for computing the exponent: the result, and the $\langle final\ sign \rangle$ are then given to `_fp_sanitizew` which checks for overflow. The exponent is computed as the largest exponent #2 or #5, incremented if there is a carry. To add the significands, we decimate the smaller number by the difference between the exponents. This is done by `_fp_add_big_i:wNww` or `_fp_add_big_ii:wNww`. We need to bring the final sign with us in the midst of the calculation to round properly at the end.

```

19061 \cs_new:Npn \__fp_add_npos_o:NnwNnw #1#2#3 ; \s__fp \__fp_chk:w 1 #4 #5
19062 {
19063   \exp_after:wN \__fp_sanitizewN
19064   \exp_after:wN #1
19065   \int_value:w \__fp_int_eval:w
19066   \if_int_compare:w #2 > #5 \exp_stop_f:
19067     #2
19068     \exp_after:wN \__fp_add_big_i_o:wNnw \int_value:w -
19069   \else:
19070     #5
19071     \exp_after:wN \__fp_add_big_ii_o:wNnw \int_value:w
19072   \fi:
19073   \__fp_int_eval:w #5 - #2 ; #1 #3;
19074 }

```

(End definition for _fp_add_npos_o:NnwNnw.)

$$\backslash_fp_add_big_i_o:wNww \langle shift \rangle ; \langle final\ sign \rangle \langle body_1 \rangle ; \langle body_2 \rangle ;$$

```

\__fp_add_big_ii_o:wNww      Used in l3fp-exo. Shift the significand of the small number, then add with \__fp_
add significand o:NnnwnnnnN.

```

```

19075 \cs_new:Npn \__fp_add_big_i_o:wNww #1; #2 #3; #4;
19076 {
19077     \__fp_decimate:nNnnnn {#1}
19078     \__fp_add_significand_o:NnnwnnnnnN
19079     #4
19080     #3
19081     #2
19082 }
19083 \cs_new:Npn \__fp_add_big_ii_o:wNww #1; #2 #3; #4;
19084 {
19085     \__fp_decimate:nNnnnn {#1}
19086     \__fp_add_significand_o:NnnwnnnnnN
19087     #3
19088     #4
19089     #2
19090 }

```

(End definition for \ fp add big i o:wNww and \ fp add big ii o:wNww.)

```

\__fp_add_significand_o:NnnwnnnnN \__fp_add_significand_o:NnnwnnnnN <rounding digit> {\langle Y'_1 \rangle} {\langle Y'_2 \rangle}
\__fp_add_significand_pack:NNNNNNN <extra-digits> ; {\langle X_1 \rangle} {\langle X_2 \rangle} {\langle X_3 \rangle} {\langle X_4 \rangle} <final sign>
\__fp_add_significand_test_o:N

```

To round properly, we must know at which digit the rounding should occur. This requires to know whether the addition produces an overall carry or not. Thus, we do the computation now and check for a carry, then go back and do the rounding. The rounding may cause a carry in very rare cases such as $0.99\dots95 \rightarrow 1.00\dots0$, but this situation always give an exact power of 10, for which it is easy to correct the result at the end.

```

19091 \cs_new:Npn \__fp_add_significand_o:NnnwnnnnN #1 #2#3 #4; #5#6#7#8
19092 {
19093   \exp_after:wN \__fp_add_significand_test_o:N
19094   \int_value:w \__fp_int_eval:w 1#5#6 + #2
19095   \exp_after:wN \__fp_add_significand_pack:NNNNNNN
19096   \int_value:w \__fp_int_eval:w 1#7#8 + #3 ; #1
19097 }
19098 \cs_new:Npn \__fp_add_significand_pack:NNNNNNN #1 #2#3#4#5#6#7
19099 {
19100   \if_meaning:w 2 #1
19101     + 1
19102   \fi:
19103   ; #2 #3 #4 #5 #6 #7 ;
19104 }
19105 \cs_new:Npn \__fp_add_significand_test_o:N #1
19106 {
19107   \if_meaning:w 2 #1
19108     \exp_after:wN \__fp_add_significand_carry_o:wwwNN
19109   \else:
19110     \exp_after:wN \__fp_add_significand_no_carry_o:wwwNN
19111   \fi:
19112 }

```

(End definition for __fp_add_significand_o:NnnwnnnnN, __fp_add_significand_pack:NNNNNNN, and __fp_add_significand_test_o:N.)

```

\__fp_add_significand_no_carry_o:wwwNN \__fp_add_significand_no_carry_o:wwwNN <8d> ; <6d> ; <2d> ; <rounding
digit> <sign>

```

If there's no carry, grab all the digits again and round. The packing function __fp_basics_pack_high:NNNNNw takes care of the case where rounding brings a carry.

```

19113 \cs_new:Npn \__fp_add_significand_no_carry_o:wwwNN
19114   #1; #2; #3#4 ; #5#6
19115 {
19116   \exp_after:wN \__fp_basics_pack_high:NNNNNw
19117   \int_value:w \__fp_int_eval:w 1 #1
19118   \exp_after:wN \__fp_basics_pack_low:NNNNNw
19119   \int_value:w \__fp_int_eval:w 1 #2 #3#4
19120   + \__fp_round:NNN #6 #4 #5
19121   \exp_after:wN ;
19122 }

```

(End definition for __fp_add_significand_no_carry_o:wwwNN.)

```

\__fp_add_significand_carry_o:wwwNN \__fp_add_significand_carry_o:wwwNN <8d> ; <6d> ; <2d> ; <rounding
digit> <sign>

```

The case where there is a carry is very similar. Rounding can even raise the first digit from 1 to 2, but we don't care.

```

19123 \cs_new:Npn \__fp_add_significand_carry_o:wwwNN
19124   #1; #2; #3#4; #5#6
19125   {
19126     + 1
19127     \exp_after:wN \__fp_basics_pack_weird_high:NNNNNNNNw
19128     \int_value:w \__fp_int_eval:w 1 1 #1
19129     \exp_after:wN \__fp_basics_pack_weird_low:NNNNw
19130     \int_value:w \__fp_int_eval:w 1 #2#3 +
19131     \exp_after:wN \__fp_round:NNN
19132     \exp_after:wN #6
19133     \exp_after:wN #3
19134     \int_value:w \__fp_round_digit:Nw #4 #5 ;
19135     \exp_after:wN ;
19136   }

```

(End definition for __fp_add_significand_carry_o:wwwNN.)

31.1.3 Absolute subtraction

```

\__fp_sub_npos_o:NnwNnw \__fp_sub_npos_o:NnwNnw <sign1> <exp1> <body1> ; \s__fp \__fp_chk:w 1
\__fp_sub_eq_o:Nnwnw <initial sign2> <exp2> <body2> ;
\__fp_sub_npos_ii_o:Nnwnw

```

Rounding properly in some modes requires to know what the sign of the result will be. Thus, we start by comparing the exponents and significands. If the numbers coincide, return zero. If the second number is larger, swap the numbers and call __fp_sub_npos_i_o:Nnwnw with the opposite of $\langle sign_1 \rangle$.

```

19137 \cs_new:Npn \__fp_sub_npos_o:NnwNnw #1#2#3; \s__fp \__fp_chk:w 1 #4#5#6;
19138   {
19139     \if_case:w \__fp_compare_npos:nwnw {#2} #3; {#5} #6; \exp_stop_f:
19140     \exp_after:wN \__fp_sub_eq_o:Nnwnw
19141     \or:
19142     \exp_after:wN \__fp_sub_npos_i_o:Nnwnw
19143     \else:
19144     \exp_after:wN \__fp_sub_npos_ii_o:Nnwnw
19145     \fi:
19146     #1 {#2} #3; {#5} #6;
19147   }
19148 \cs_new:Npn \__fp_sub_eq_o:Nnwnw #1#2; #3; { \exp_after:wN \c_zero_fp }
19149 \cs_new:Npn \__fp_sub_npos_ii_o:Nnwnw #1 #2; #3;
19150   {
19151     \exp_after:wN \__fp_sub_npos_i_o:Nnwnw
19152     \int_value:w \__fp_neg_sign:N #1
19153     #3; #2;
19154   }

```

(End definition for __fp_sub_npos_o:NnwNnw, __fp_sub_eq_o:Nnwnw, and __fp_sub_npos_ii_o:Nnwnw.)

__fp_sub_npos_i_o:Nnwnw After the computation is done, __fp_sanitize:Nw checks for overflow/underflow. It expects the $\langle final\ sign \rangle$ and the $\langle exponent \rangle$ (delimited by ;). Start an integer expression for the exponent, which starts with the exponent of the largest number, and may be decreased if the two numbers are very close. If the two numbers have the same exponent, call the `near` auxiliary. Otherwise, decimate y , then call the `far` auxiliary to evaluate

the difference between the two significands. Note that we decimate by 1 less than one could expect.

```

19155 \cs_new:Npn \__fp_sub_npos_i_o:Nnwnw #1 #2#3; #4#5;
19156 {
19157   \exp_after:wN \__fp_sanitizew
19158   \exp_after:wN #1
19159   \int_value:w \__fp_int_eval:w
19160   #2
19161   \if_int_compare:w #2 = #4 \exp_stop_f:
19162     \exp_after:wN \__fp_sub_back_near_o:nnnnnnnnN
19163   \else:
19164     \exp_after:wN \__fp_decimate:nNnnnn \exp_after:wN
19165     { \int_value:w \__fp_int_eval:w #2 - #4 - 1 \exp_after:wN }
19166     \exp_after:wN \__fp_sub_back_far_o:NnnwnnnnnN
19167   \fi:
19168   #5
19169   #3
19170   #1
19171 }

```

(End definition for __fp_sub_npos_i_o:Nnwnw.)

```

\__fp_sub_back_near_o:nnnnnnnnN      \__fp_sub_back_near_o:nnnnnnnnN {<Y1>} {<Y2>} {<Y3>} {<Y4>} {<X1>}
\__fp_sub_back_near_pack:NNNNNNw     {<X2>} {<X3>} {<X4>} {<final sign>}
\__fp_sub_back_near_after:wNNNNw

```

In this case, the subtraction is exact, so we discard the *<final sign>* #9. The very large shifts of 10^9 and $1.1 \cdot 10^9$ are unnecessary here, but allow the auxiliaries to be reused later. Each integer expression produces a 10 digit result. If the resulting 16 digits start with a 0, then we need to shift the group, padding with trailing zeros.

```

19172 \cs_new:Npn \__fp_sub_back_near_o:nnnnnnnnN #1#2#3#4 #5#6#7#8 #9
19173 {
19174   \exp_after:wN \__fp_sub_back_near_after:wNNNNw
19175   \int_value:w \__fp_int_eval:w 10#5#6 - #1#2 - 11
19176   \exp_after:wN \__fp_sub_back_near_pack:NNNNNNw
19177   \int_value:w \__fp_int_eval:w 11#7#8 - #3#4 \exp_after:wN ;
19178 }
19179 \cs_new:Npn \__fp_sub_back_near_pack:NNNNNNw #1#2#3#4#5#6#7 ;
19180 { + #1#2 ; {#3#4#5#6} {#7} ; }
19181 \cs_new:Npn \__fp_sub_back_near_after:wNNNNw 10 #1#2#3#4 #5 ;
19182 {
19183   \if_meaning:w 0 #1
19184     \exp_after:wN \__fp_sub_back_shift:wnnnn
19185   \fi:
19186   ; {#1#2#3#4} {#5}
19187 }

```

(End definition for __fp_sub_back_near_o:nnnnnnnnN, __fp_sub_back_near_pack:NNNNNNw, and __fp_sub_back_near_after:wNNNNw.)

```

\__fp_sub_back_shift:wnnnn      \__fp_sub_back_shift:wnnnn ; {<Z1>} {<Z2>} {<Z3>} {<Z4>} ;
\__fp_sub_back_shift_ii:ww      This function is called with <Z1> ≤ 999. Act with \number to trim leading zeros from
\__fp_sub_back_shift_iii:NNNNNNw <Z1> <Z2> (we don't do all four blocks at once, since non-zero blocks would then overflow
\__fp_sub_back_shift_iv:nnnnw   TEX's integers). If the first two blocks are zero, the auxiliary receives an empty #1 and

```

trims #2#30 from leading zeros, yielding a total shift between 7 and 16 to the exponent. Otherwise we get the shift from #1 alone, yielding a result between 1 and 6. Once the

exponent is taken care of, trim leading zeros from #1#2#3 (when #1 is empty, the space before #2#3 is ignored), get four blocks of 4 digits and finally clean up. Trailing zeros are added so that digits can be grabbed safely.

```

19188 \cs_new:Npn \__fp_sub_back_shift:wnnnn ; #1#2
19189 {
19190   \exp_after:wN \__fp_sub_back_shift_ii:ww
19191   \int_value:w #1 #2 0 ;
19192 }
19193 \cs_new:Npn \__fp_sub_back_shift_ii:ww #1 0 ; #2#3 ;
19194 {
19195   \if_meaning:w @ #1 @
19196   - 7
19197   - \exp_after:wN \use_i:nnn
19198     \exp_after:wN \__fp_sub_back_shift_iii:NNNNNNNNw
19199     \int_value:w #2#3 0 ~ 123456789;
19200   \else:
19201     - \__fp_sub_back_shift_iii:NNNNNNNNw #1 123456789;
19202   \fi:
19203   \exp_after:wN \__fp_pack_twice_four:wNNNNNNNN
19204   \exp_after:wN \__fp_pack_twice_four:wNNNNNNNN
19205   \exp_after:wN \__fp_sub_back_shift_iv:nnnnw
19206   \exp_after:wN ;
19207   \int_value:w
19208   #1 ~ #2#3 0 ~ 0000 0000 0000 000 ;
19209 }
19210 \cs_new:Npn \__fp_sub_back_shift_iii:NNNNNNNNw #1#2#3#4#5#6#7#8#9; {#8}
19211 \cs_new:Npn \__fp_sub_back_shift_iv:nnnnw #1 ; #2 ; { ; #1 ; }

```

(End definition for __fp_sub_back_shift:wnnnn and others.)

__fp_sub_back_far_o:NnnwnnnnN $\langle \text{rounding} \rangle \{ \langle Y'_1 \rangle \} \{ \langle Y'_2 \rangle \}$
 $\langle \text{extra-digits} \rangle ; \{ \langle X_1 \rangle \} \{ \langle X_2 \rangle \} \{ \langle X_3 \rangle \} \{ \langle X_4 \rangle \} \langle \text{final sign} \rangle$

If the difference is greater than $10^{\langle expo_x \rangle}$, call the `very_far` auxiliary. If the result is less than $10^{\langle expo_x \rangle}$, call the `not_far` auxiliary. If it is too close a call to know yet, namely if $1 \langle Y'_1 \rangle \langle Y'_2 \rangle = \langle X_1 \rangle \langle X_2 \rangle \langle X_3 \rangle \langle X_4 \rangle 0$, then call the `quite_far` auxiliary. We use the odd combination of space and semi-colon delimiters to allow the `not_far` auxiliary to grab each piece individually, the `very_far` auxiliary to use `__fp_pack_eight:wNNNNNNNN`, and the `quite_far` to ignore the significands easily (using the `; delimiter`).

```

19212 \cs_new:Npn \__fp_sub_back_far_o:NnnwnnnnN #1 #2#3 #4; #5#6#7#8
19213 {
19214   \if_case:w
19215     \if_int_compare:w 1 #2 = #5#6 \use_i:nnnn #7 \exp_stop_f:
19216     \if_int_compare:w #3 = \use_none:n #7#8 0 \exp_stop_f:
19217     0
19218     \else:
19219       \if_int_compare:w #3 > \use_none:n #7#8 0 - \fi: 1
19220       \fi:
19221     \else:
19222       \if_int_compare:w 1 #2 > #5#6 \use_i:nnnn #7 - \fi: 1
19223       \fi:
19224     \exp_stop_f:
19225     \exp_after:wN \__fp_sub_back_quite_far_o:wwNN
19226   \or: \exp_after:wN \__fp_sub_back_very_far_o:wwwNN

```

```

19227     \else: \exp_after:wN \_fp_sub_back_not_far_o:wwwNN
19228     \fi:
19229     #2 ~ #3 ; #5 #6 ~ #7 #8 ; #1
19230 }

```

(End definition for _fp_sub_back_far_o:NnnwnnnN.)

_fp_sub_back_quite_far_o:wwNN
_fp_sub_back_quite_far_ii:NN

The easiest case is when $x - y$ is extremely close to a power of 10, namely the first digit of x is 1, and all others vanish when subtracting y . Then the *rounding* #3 and the *final sign* #4 control whether we get 1 or 0.9999999999999999. In the usual round-to-nearest mode, we get 1 whenever the *rounding* digit is less than or equal to 5 (remember that the *rounding* digit is only equal to 5 if there was no further non-zero digit).

```

19231 \cs_new:Npn \_fp_sub_back_quite_far_o:wwNN #1; #2; #3#4
19232 {
19233     \exp_after:wN \_fp_sub_back_quite_far_ii:NN
19234     \exp_after:wN #3
19235     \exp_after:wN #4
19236 }
19237 \cs_new:Npn \_fp_sub_back_quite_far_ii:NN #1#2
19238 {
19239     \if_case:w \_fp_round_neg:NNN #2 0 #1
19240     \exp_after:wN \use_i:nn
19241     \else:
19242     \exp_after:wN \use_ii:nn
19243     \fi:
19244     { ; {1000} {0000} {0000} {0000} ; }
19245     { - 1 ; {9999} {9999} {9999} {9999} ; }
19246 }

```

(End definition for _fp_sub_back_quite_far_o:wwNN and _fp_sub_back_quite_far_ii:NN.)

_fp_sub_back_not_far_o:wwwNN

In the present case, x and y have different exponents, but y is large enough that $x - y$ has a smaller exponent than x . Decrement the exponent (with -1). Then proceed in a way similar to the **near** auxiliaries seen earlier, but multiplying x by 10 (#30 and #40 below), and with the added quirk that the *rounding* digit has to be taken into account. Namely, we may have to decrease the result by one unit if _fp_round_neg:NNN returns 1. This function expects the *final sign* #6, the last digit of 1100000000+#40-#2, and the *rounding* digit. Instead of redoing the computation for the second argument, we note that _fp_round_neg:NNN only cares about its parity, which is identical to that of the last digit of #2.

```

19247 \cs_new:Npn \_fp_sub_back_not_far_o:wwwNN #1 ~ #2; #3 ~ #4; #5#6
19248 {
19249     - 1
19250     \exp_after:wN \_fp_sub_back_near_after:wNNNNw
19251     \int_value:w \_fp_int_eval:w 1#30 - #1 - 11
19252     \exp_after:wN \_fp_sub_back_near_pack:NNNNNNw
19253     \int_value:w \_fp_int_eval:w 11 0000 0000 + #40 - #2
19254     - \exp_after:wN \_fp_round_neg:NNN
19255     \exp_after:wN #6
19256     \use_none:nnnnnn #2 #5
19257     \exp_after:wN ;
19258 }

```

(End definition for _fp_sub_back_not_far_o:wwwNN.)

`_fp_sub_back_very_far_o:wwwNN`
`_fp_sub_back_very_far_ii_o:nnNwwNN`

The case where $x - y$ and x have the same exponent is a bit more tricky, mostly because it cannot reuse the same auxiliaries. Shift the y significand by adding a leading 0. Then the logic is similar to the `not_far` functions above. Rounding is a bit more complicated: we have two *rounding* digits #3 and #6 (from the decimation, and from the new shift) to take into account, and getting the parity of the main result requires a computation. The first `\int_value:w` triggers the second one because the number is unfinished; we can thus not use 0 in place of 2 there.

```

19259 \cs_new:Npn \_fp_sub_back_very_far_o:wwwNN #1#2#3#4#5#6#7
19260 {
19261   \_fp_pack_eight:wNNNNNNNN
19262   \_fp_sub_back_very_far_ii_o:nnNwwNN
19263   { 0 #1#2#3 #4#5#6#7 }
19264   ;
19265 }
19266 \cs_new:Npn \_fp_sub_back_very_far_ii_o:nnNwwNN #1#2 ; #3 ; #4 ~ #5; #6#7
19267 {
19268   \exp_after:wN \_fp_basics_pack_high:NNNNw
19269   \int_value:w \_fp_int_eval:w 1#4 - #1 - 1
19270   \exp_after:wN \_fp_basics_pack_low:NNNNw
19271   \int_value:w \_fp_int_eval:w 2#5 - #2
19272   - \exp_after:wN \_fp_round_neg:NNN
19273   \exp_after:wN #7
19274   \int_value:w
19275   \if_int_odd:w \_fp_int_eval:w #5 - #2 \_fp_int_eval_end:
19276   1 \else: 2 \fi:
19277   \int_value:w \_fp_round_digit:Nw #3 #6 ;
19278   \exp_after:wN ;
19279 }

```

(End definition for `_fp_sub_back_very_far_o:wwwNN` and `_fp_sub_back_very_far_ii_o:nnNwwNN`.)

31.2 Multiplication

31.2.1 Signs, and special numbers

`_fp*_o:ww`

We go through an auxiliary, which is common with `_fp/_o:ww`. The first argument is the operation, used for the invalid operation exception. The second is inserted in a formula to dispatch cases slightly differently between multiplication and division. The third is the operation for normal floating points. The fourth is there for extra cases needed in `_fp/_o:ww`.

```

19280 \cs_new:cpn { \_fp*_o:ww }
19281 {
19282   \_fp_mul_cases_o:NnNww
19283   *
19284   { - 2 + }
19285   \_fp_mul_npos_o:Nww
19286   { }
19287 }

```

(End definition for `_fp*_o:ww`.)

`_fp_mul_cases_o:nNnnww`

Split into 10 cases (12 for division). If both numbers are normal, go to case 0 (same sign) or case 1 (opposite signs): in both cases, call `_fp_mul_npos_o:Nww` to do the work. If

the first operand is `nan`, go to case 2, in which the second operand is discarded; if the second operand is `nan`, go to case 3, in which the first operand is discarded (note the weird interaction with the final test on signs). Then we separate the case where the first number is normal and the second is zero: this goes to cases 4 and 5 for multiplication, 10 and 11 for division. Otherwise, we do a computation which dispatches the products $0 \times 0 = 0 \times 1 = 1 \times 0 = 0$ to case 4 or 5 depending on the combined sign, the products $0 \times \infty$ and $\infty \times 0$ to case 6 or 7 (invalid operation), and the products $1 \times \infty = \infty \times 1 = \infty \times \infty = \infty$ to cases 8 and 9. Note that the code for these two cases (which return $\pm\infty$) is inserted as argument #4, because it differs in the case of divisions.

```

19288 \cs_new:Npn \__fp_mul_cases_o:NnNnw
19289 #1#2#3#4 \s__fp \__fp_chk:w #5#6#7; \s__fp \__fp_chk:w #8#9
19290 {
19291   \if_case:w \__fp_int_eval:w
19292     \if_int_compare:w #5 #8 = 11 ~
19293       1
19294     \else:
19295       \if_meaning:w 3 #8
19296       3
19297     \else:
19298       \if_meaning:w 3 #5
19299       2
19300     \else:
19301       \if_int_compare:w #5 #8 = 10 ~
19302       9 #2 - 2
19303     \else:
19304       (#5 #2 #8) / 2 * 2 + 7
19305     \fi:
19306   \fi:
19307   \fi:
19308   \fi:
19309   \if_meaning:w #6 #9 - 1 \fi:
19310   \__fp_int_eval_end:
19311   \__fp_case_use:nw { #3 0 }
19312   \or: \__fp_case_use:nw { #3 2 }
19313   \or: \__fp_case_return_i_o:ww
19314   \or: \__fp_case_return_ii_o:ww
19315   \or: \__fp_case_return_o:Nww \c_zero_fp
19316   \or: \__fp_case_return_o:Nww \c_minus_zero_fp
19317   \or: \__fp_case_use:nw { \__fp_invalid_operation_o:Nww #1 }
19318   \or: \__fp_case_use:nw { \__fp_invalid_operation_o:Nww #1 }
19319   \or: \__fp_case_return_o:Nww \c_inf_fp
19320   \or: \__fp_case_return_o:Nww \c_minus_inf_fp
19321   #4
19322   \fi:
19323   \s__fp \__fp_chk:w #5 #6 #7;
19324   \s__fp \__fp_chk:w #8 #9
19325 }

```

(End definition for `__fp_mul_cases_o:nNnnnw`.)

31.2.2 Absolute multiplication

In this subsection, we perform the multiplication of two positive normal numbers.

```

\__fp_mul_npos_o:Nww \__fp_mul_npos_o:Nww <final sign> \s__fp \__fp_chk:w 1 <sign1> {<exp1>}
<body1> ; \s__fp \__fp_chk:w 1 <sign2> {<exp2>} <body2> ;

```

After the computation, __fp_sanitize:Nw checks for overflow or underflow. As we did for addition, __fp_int_eval:w computes the exponent, catching any shift coming from the computation in the significand. The <final sign> is needed to do the rounding properly in the significand computation. We setup the post-expansion here, triggered by __fp_mul_significand_o:nnnnNnnnn.

This is also used in l3fp-convert.

```

19326 \cs_new:Npn \__fp_mul_npos_o:Nww
19327 #1 \s__fp \__fp_chk:w #2 #3 #4 #5 ; \s__fp \__fp_chk:w #6 #7 #8 #9 ;
19328 {
19329   \exp_after:wN \__fp_sanitize:Nw
19330   \exp_after:wN #1
19331   \int_value:w \__fp_int_eval:w
19332     #4 + #8
19333   \__fp_mul_significand_o:nnnnNnnnn #5 #1 #9
19334 }

```

(End definition for __fp_mul_npos_o:Nww.)

```

\__fp_mul_significand_o:nnnnNnnnn \__fp_mul_significand_o:nnnnNnnnn {<X1>} {<X2>} {<X3>} {<X4>} <sign>
\__fp_mul_significand_drop:NNNNNw {<Y1>} {<Y2>} {<Y3>} {<Y4>}
\__fp_mul_significand_keep:NNNNNw

```

Note the three semicolons at the end of the definition. One is for the last __fp_mul_significand_drop:NNNNNw; one is for __fp_round_digit:Nw later on; and one, preceded by \exp_after:wN, which is correctly expanded (within an __fp_int_eval:w), is used by __fp_basics_pack_low:NNNNNw.

The product of two 16 digit integers has 31 or 32 digits, but it is impossible to know which one before computing. The place where we round depends on that number of digits, and may depend on all digits until the last in some rare cases. The approach is thus to compute the 5 first blocks of 4 digits (the first one is between 100 and 9999 inclusive), and a compact version of the remaining 3 blocks. Afterwards, the number of digits is known, and we can do the rounding within yet another set of __fp_int_eval:w.

```

19335 \cs_new:Npn \__fp_mul_significand_o:nnnnNnnnn #1#2#3#4 #5 #6#7#8#9
19336 {
19337   \exp_after:wN \__fp_mul_significand_test_f:NNN
19338   \exp_after:wN #5
19339   \int_value:w \__fp_int_eval:w 99990000 + #1*#6 +
19340   \exp_after:wN \__fp_mul_significand_keep:NNNNNw
19341   \int_value:w \__fp_int_eval:w 99990000 + #1*#7 + #2*#6 +
19342   \exp_after:wN \__fp_mul_significand_keep:NNNNNw
19343   \int_value:w \__fp_int_eval:w 99990000 + #1*#8 + #2*#7 + #3*#6 +
19344   \exp_after:wN \__fp_mul_significand_drop:NNNNNw
19345   \int_value:w \__fp_int_eval:w 99990000 + #1*#9 + #2*#8 +
19346   #3*#7 + #4*#6 +
19347   \exp_after:wN \__fp_mul_significand_drop:NNNNNw
19348   \int_value:w \__fp_int_eval:w 99990000 + #2*#9 + #3*#8 +
19349   #4*#7 +
19350   \exp_after:wN \__fp_mul_significand_drop:NNNNNw
19351   \int_value:w \__fp_int_eval:w 99990000 + #3*#9 + #4*#8 +
19352   \exp_after:wN \__fp_mul_significand_drop:NNNNNw
19353   \int_value:w \__fp_int_eval:w 100000000 + #4*#9 ;
19354 ; \exp_after:wN ;

```

```

19355     }
19356 \cs_new:Npn \__fp_mul_significand_drop:NNNNNw #1#2#3#4#5 #6;
19357 { #1#2#3#4#5 ; + #6 }
19358 \cs_new:Npn \__fp_mul_significand_keep:NNNNNw #1#2#3#4#5 #6;
19359 { #1#2#3#4#5 ; #6 ; }

```

(End definition for __fp_mul_significand_o:nnnnNnnnn, __fp_mul_significand_drop:NNNNNw, and __fp_mul_significand_keep:NNNNNw.)

```

\__fp_mul_significand_test_f:NNN \__fp_mul_significand_test_f:NNN <sign> 1 <digits 1-8> ; <digits 9-12> ;
<digits 13-16> ; + <digits 17-20> + <digits 21-24> + <digits 25-28> + <digits
29-32> ; \exp_after:wN ;

```

If the $\langle \text{digit } 1 \rangle$ is non-zero, then for rounding we only care about the digits 16 and 17, and whether further digits are zero or not (check for exact ties). On the other hand, if $\langle \text{digit } 1 \rangle$ is zero, we care about digits 17 and 18, and whether further digits are zero.

```

19360 \cs_new:Npn \__fp_mul_significand_test_f:NNN #1 #2 #3
19361 {
19362   \if_meaning:w 0 #3
19363     \exp_after:wN \__fp_mul_significand_small_f:NNwwwN
19364   \else:
19365     \exp_after:wN \__fp_mul_significand_large_f:NwwNNNN
19366   \fi:
19367   #1 #3
19368 }

```

(End definition for __fp_mul_significand_test_f:NNN.)

__fp_mul_significand_large_f:NwwNNNN In this branch, $\langle \text{digit } 1 \rangle$ is non-zero. The result is thus $\langle \text{digits } 1-16 \rangle$, plus some rounding which depends on the digits 16, 17, and whether all subsequent digits are zero or not. Here, __fp_round_digit:Nw takes digits 17 and further (as an integer expression), and replaces it by a $\langle \text{rounding digit} \rangle$, suitable for __fp_round:NNN.

```

19369 \cs_new:Npn \__fp_mul_significand_large_f:NwwNNNN #1 #2; #3; #4#5#6#7; +
19370 {
19371   \exp_after:wN \__fp_basics_pack_high:NNNNNw
19372   \int_value:w \__fp_int_eval:w 1#2
19373   \exp_after:wN \__fp_basics_pack_low:NNNNNw
19374   \int_value:w \__fp_int_eval:w 1#3#4#5#6#7
19375   + \exp_after:wN \__fp_round:NNN
19376   \exp_after:wN #1
19377   \exp_after:wN #7
19378   \int_value:w \__fp_round_digit:Nw
19379 }

```

(End definition for __fp_mul_significand_large_f:NwwNNNN.)

__fp_mul_significand_small_f:NNwwwN In this branch, $\langle \text{digit } 1 \rangle$ is zero. Our result is thus $\langle \text{digits } 2-17 \rangle$, plus some rounding which depends on the digits 17, 18, and whether all subsequent digits are zero or not. The 8 digits 1#3 are followed, after expansion of the small_pack auxiliary, by the next digit, to form a 9 digit number.

```

19380 \cs_new:Npn \__fp_mul_significand_small_f:NNwwwN #1 #2#3; #4#5; #6; + #7
19381 {
19382   - 1
19383   \exp_after:wN \__fp_basics_pack_high:NNNNNw
19384   \int_value:w \__fp_int_eval:w 1#3#4

```

```

19385     \exp_after:wN \__fp_basics_pack_low:NNNNw
19386     \int_value:w \__fp_int_eval:w 1#5#6#7
19387     + \exp_after:wN \__fp_round:NNN
19388       \exp_after:wN #1
19389       \exp_after:wN #7
19390       \int_value:w \__fp_round_digit:Nw
19391   }

```

(End definition for __fp_mul_significand_small_f:NNwwN.)

31.3 Division

31.3.1 Signs, and special numbers

Time is now ripe to tackle the hardest of the four elementary operations: division.

__fp/_o:ww Filtering special floating point is very similar to what we did for multiplications, with a few variations. Invalid operation exceptions display / rather than *. In the formula for dispatch, we replace - 2 + by -. The case of normal numbers is treated using __fp_div_npos_o:Nww rather than __fp_mul_npos_o:Nww. There are two additional cases: if the first operand is normal and the second is a zero, then the division by zero exception is raised: cases 10 and 11 of the \if_case:w construction in __fp_mul_cases_o:NnNww are provided as the fourth argument here.

```

19392 \cs_new:cpn { __fp/_o:ww }
19393 {
19394     \__fp_mul_cases_o:NnNww
19395     /
19396     { - }
19397     \__fp_div_npos_o:Nww
19398     {
19399         \or:
19400             \__fp_case_use:nw
19401             { \__fp_division_by_zero_o:NNww \c_inf_fp / }
19402         \or:
19403             \__fp_case_use:nw
19404             { \__fp_division_by_zero_o:NNww \c_minus_inf_fp / }
19405     }
19406 }

```

(End definition for __fp/_o:ww.)

```

\__fp_div_npos_o:Nww \__fp_div_npos_o:Nww <final sign> \s__fp \__fp_chk:w 1 <sign_A> {\<exp A>}
{\<A_1>} {\<A_2>} {\<A_3>} {\<A_4>} ; \s__fp \__fp_chk:w 1 <sign_Z> {\<exp Z>}
{\<Z_1>} {\<Z_2>} {\<Z_3>} {\<Z_4>} ;

```

We want to compute A/Z . As for multiplication, __fp_sanitize:Nw checks for overflow or underflow; we provide it with the $\langle final\ sign \rangle$, and an integer expression in which we compute the exponent. We set up the arguments of __fp_div_significand_i_o:wnnw, namely an integer $\langle y \rangle$ obtained by adding 1 to the first 5 digits of Z (explanation given soon below), then the four $\{\langle A_i \rangle\}$, then the four $\{\langle Z_i \rangle\}$, a semi-colon, and the $\langle final\ sign \rangle$, used for rounding at the end.

```

19407 \cs_new:Npn \__fp_div_npos_o:Nww
19408     #1 \s__fp \__fp_chk:w 1 #2 #3 #4 ; \s__fp \__fp_chk:w 1 #5 #6 #7#8#9;
19409 {

```

```

19410      \exp_after:wN \__fp_sanitizew
19411      \exp_after:wN #1
19412      \int_value:w \__fp_int_eval:w
19413      #3 - #6
19414      \exp_after:wN \__fp_div_significand_i_o:wnnw
19415      \int_value:w \__fp_int_eval:w #7 \use_i:nnnn #8 + 1 ;
19416      #4
19417      {#7}{#8}#9 ;
19418      #1
19419  }

```

(End definition for `__fp_div_npos_o:Nww`.)

31.3.2 Work plan

In this subsection, we explain how to avoid overflowing \TeX 's integers when performing the division of two positive normal numbers.

We are given two numbers, $A = 0.A_1A_2A_3A_4$ and $Z = 0.Z_1Z_2Z_3Z_4$, in blocks of 4 digits, and we know that the first digits of A_1 and of Z_1 are non-zero. To compute A/Z , we proceed as follows.

- Find an integer $Q_A \simeq 10^4 A/Z$.
- Replace A by $B = 10^4 A - Q_A Z$.
- Find an integer $Q_B \simeq 10^4 B/Z$.
- Replace B by $C = 10^4 B - Q_B Z$.
- Find an integer $Q_C \simeq 10^4 C/Z$.
- Replace C by $D = 10^4 C - Q_C Z$.
- Find an integer $Q_D \simeq 10^4 D/Z$.
- Consider $E = 10^4 D - Q_D Z$, and ensure correct rounding.

The result is then $Q = 10^{-4}Q_A + 10^{-8}Q_B + 10^{-12}Q_C + 10^{-16}Q_D + \text{rounding}$. Since the Q_i are integers, B , C , D , and E are all exact multiples of 10^{-16} , in other words, computing with 16 digits after the decimal separator yields exact results. The problem is the risk of overflow: in general B , C , D , and E may be greater than 1.

Unfortunately, things are not as easy as they seem. In particular, we want all intermediate steps to be positive, since negative results would require extra calculations at the end. This requires that $Q_A \leq 10^4 A/Z$ etc. A reasonable attempt would be to define Q_A as

$$\text{\int_eval:n} \left\{ \frac{A_1 A_2}{Z_1 + 1} - 1 \right\} \leq 10^4 \frac{A}{Z}$$

Subtracting 1 at the end takes care of the fact that $\varepsilon\text{-}\text{\TeX}$'s `__fp_int_eval:w` rounds divisions instead of truncating (really, $1/2$ would be sufficient, but we work with integers). We add 1 to Z_1 because $Z_1 \leq 10^4 Z < Z_1 + 1$ and we need Q_A to be an underestimate. However, we are now underestimating Q_A too much: it can be wrong by up to 100, for instance when $Z = 0.1$ and $A \simeq 1$. Then B could take values up to 10 (maybe more), and a few steps down the line, we would run into arithmetic overflow, since \TeX can only handle integers less than roughly $2 \cdot 10^9$.

A better formula is to take

$$Q_A = \backslash\text{int_eval:n}\left\{\frac{10 \cdot A_1 A_2}{\lfloor 10^{-3} \cdot Z_1 Z_2 \rfloor + 1} - 1\right\}.$$

This is always less than $10^9 A / (10^5 Z)$, as we wanted. In words, we take the 5 first digits of Z into account, and the 8 first digits of A , using 0 as a 9-th digit rather than the true digit for efficiency reasons. We shall prove that using this formula to define all the Q_i avoids any overflow. For convenience, let us denote

$$y = \lfloor 10^{-3} \cdot Z_1 Z_2 \rfloor + 1,$$

so that, taking into account the fact that $\varepsilon\text{-TeX}$ rounds ties away from zero,

$$\begin{aligned} Q_A &= \left\lfloor \frac{A_1 A_2 0}{y} - \frac{1}{2} \right\rfloor \\ &> \frac{A_1 A_2 0}{y} - \frac{3}{2}. \end{aligned}$$

Note that $10^4 < y \leq 10^5$, and $999 \leq Q_A \leq 99989$. Also note that this formula does not cause an overflow as long as $A < (2^{31} - 1)/10^9 \simeq 2.147 \dots$, since the numerator involves an integer slightly smaller than $10^9 A$.

Let us bound B :

$$\begin{aligned} 10^5 B &= A_1 A_2 0 + 10 \cdot 0.A_3 A_4 - 10 \cdot Z_1.Z_2 Z_3 Z_4 \cdot Q_A \\ &< A_1 A_2 0 \cdot \left(1 - 10 \cdot \frac{Z_1.Z_2 Z_3 Z_4}{y}\right) + \frac{3}{2} \cdot 10 \cdot Z_1.Z_2 Z_3 Z_4 + 10 \\ &\leq \frac{A_1 A_2 0 \cdot (y - 10 \cdot Z_1.Z_2 Z_3 Z_4)}{y} + \frac{3}{2} y + 10 \\ &\leq \frac{A_1 A_2 0 \cdot 1}{y} + \frac{3}{2} y + 10 \leq \frac{10^9 A}{y} + 1.6 \cdot y. \end{aligned}$$

At the last step, we hide 10 into the second term for later convenience. The same reasoning yields

$$\begin{aligned} 10^5 B &< 10^9 A / y + 1.6y, \\ 10^5 C &< 10^9 B / y + 1.6y, \\ 10^5 D &< 10^9 C / y + 1.6y, \\ 10^5 E &< 10^9 D / y + 1.6y. \end{aligned}$$

The goal is now to prove that none of B , C , D , and E can go beyond $(2^{31} - 1)/10^9 = 2.147 \dots$.

Combining the various inequalities together with $A < 1$, we get

$$\begin{aligned} 10^5 B &< 10^9 / y + 1.6y, \\ 10^5 C &< 10^{13} / y^2 + 1.6(y + 10^4), \\ 10^5 D &< 10^{17} / y^3 + 1.6(y + 10^4 + 10^8 / y), \\ 10^5 E &< 10^{21} / y^4 + 1.6(y + 10^4 + 10^8 / y + 10^{12} / y^2). \end{aligned}$$

All of those bounds are convex functions of y (since every power of y involved is convex, and the coefficients are positive), and thus maximal at one of the end-points of the allowed range $10^4 < y \leq 10^5$. Thus,

$$\begin{aligned} 10^5 B &< \max(1.16 \cdot 10^5, 1.7 \cdot 10^5), \\ 10^5 C &< \max(1.32 \cdot 10^5, 1.77 \cdot 10^5), \\ 10^5 D &< \max(1.48 \cdot 10^5, 1.777 \cdot 10^5), \\ 10^5 E &< \max(1.64 \cdot 10^5, 1.7777 \cdot 10^5). \end{aligned}$$

All of those bounds are less than $2.147 \cdot 10^5$, and we are thus within $\text{T}_{\text{E}}\text{X}$'s bounds in all cases!

We later need to have a bound on the Q_i . Their definitions imply that $Q_A < 10^9 A/y - 1/2 < 10^5 A$ and similarly for the other Q_i . Thus, all of them are less than 177770.

The last step is to ensure correct rounding. We have

$$A/Z = \sum_{i=1}^4 (10^{-4i} Q_i) + 10^{-16} E/Z$$

exactly. Furthermore, we know that the result is in $[0.1, 10)$, hence will be rounded to a multiple of 10^{-16} or of 10^{-15} , so we only need to know the integer part of E/Z , and a “rounding” digit encoding the rest. Equivalently, we need to find the integer part of $2E/Z$, and determine whether it was an exact integer or not (this serves to detect ties). Since

$$\frac{2E}{Z} = 2 \frac{10^5 E}{10^5 Z} \leq 2 \frac{10^5 E}{10^4} < 36,$$

this integer part is between 0 and 35 inclusive. We let $\varepsilon\text{-T}_{\text{E}}\text{X}$ round

$$P = \backslash\text{int_eval:n} \left\{ \frac{2 \cdot E_1 E_2}{Z_1 Z_2} \right\},$$

which differs from $2E/Z$ by at most

$$\frac{1}{2} + 2 \left| \frac{E}{Z} - \frac{E}{10^{-8} Z_1 Z_2} \right| + 2 \left| \frac{10^8 E - E_1 E_2}{Z_1 Z_2} \right| < 1,$$

($1/2$ comes from $\varepsilon\text{-T}_{\text{E}}\text{X}$'s rounding) because each absolute value is less than 10^{-7} . Thus P is either the correct integer part, or is off by 1; furthermore, if $2E/Z$ is an integer, $P = 2E/Z$. We will check the sign of $2E - PZ$. If it is negative, then $E/Z \in ((P-1)/2, P/2)$. If it is zero, then $E/Z = P/2$. If it is positive, then $E/Z \in (P/2, (P+1)/2)$. In each case, we know how to round to an integer, depending on the parity of P , and the rounding mode.

31.3.3 Implementing the significand division

`_fp_div_significand_i_o:wnnw`

`_fp_div_significand_i_o:wnnw <y> ; {<A1>} {<A2>} {<A3>} {<A4>}
{<Z1>} {<Z2>} {<Z3>} {<Z4>} ; <sign>`

Compute $10^6 + Q_A$ (a 7 digit number thanks to the shift), unbrace $\langle A_1 \rangle$ and $\langle A_2 \rangle$, and prepare the $\langle continuation \rangle$ arguments for 4 consecutive calls to `_fp_div_significand_calc:wnnnnnnn`. Each of these calls needs $\langle y \rangle$ ($\#1$), and it turns out that

we need post-expansion there, hence the `\int_value:w`. Here, `#4` is six brace groups, which give the six first n-type arguments of the `calc` function.

```

19420 \cs_new:Npn \__fp_div_significand_i_o:wnnw #1 ; #2#3 #4 ;
19421 {
19422   \exp_after:wN \__fp_div_significand_test_o:w
19423   \int_value:w \__fp_int_eval:w
19424   \exp_after:wN \__fp_div_significand_calc:wnnnnnnn
19425   \int_value:w \__fp_int_eval:w 999999 + #2 #3 0 / #1 ;
19426   #2 #3 ;
19427   #4
19428   { \exp_after:wN \__fp_div_significand_ii:wN \int_value:w #1 }
19429   { \exp_after:wN \__fp_div_significand_ii:wN \int_value:w #1 }
19430   { \exp_after:wN \__fp_div_significand_ii:wN \int_value:w #1 }
19431   { \exp_after:wN \__fp_div_significand_iii:wnnnnnn \int_value:w #1 }
19432 }

```

(End definition for `__fp_div_significand_i_o:wnnw`.)

```

\__fp_div_significand_calc:wnnnnnnn \__fp_div_significand_calc:wnnnnnnn <106 + QA> ; <A1> <A2> ; {<A3>}
\__fp_div_significand_calc_i:wnnnnnnn {<A4>} {<Z1>} {<Z2>} {<Z3>} {<Z4>} {<continuation>}
\__fp_div_significand_calc_ii:wnnnnnnn expands to
<106 + QA> <continuation> ; <B1> <B2> ; {<B3>} {<B4>} {<Z1>} {<Z2>} {<Z3>}
{<Z4>}

```

where $B = 10^4 A - Q_A \cdot Z$. This function is also used to compute C , D , E (with the input shifted accordingly), and is used in `l3fp-expo`.

We know that $0 < Q_A < 1.8 \cdot 10^5$, so the product of Q_A with each Z_i is within $\text{T}_{\text{E}}\text{X}$'s bounds. However, it is a little bit too large for our purposes: we would not be able to use the usual trick of adding a large power of 10 to ensure that the number of digits is fixed.

The bound on Q_A , implies that $10^6 + Q_A$ starts with the digit 1, followed by 0 or 1. We test, and call different auxiliaries for the two cases. An earlier implementation did the tests within the computation, but since we added a `<continuation>`, this is not possible because the macro has 9 parameters.

The result we want is then (the overall power of 10 is arbitrary):

$$\begin{aligned}
& 10^{-4}(\#2 - \#1 \cdot \#5 - 10 \cdot \langle i \rangle \cdot \#5\#6) + 10^{-8}(\#3 - \#1 \cdot \#6 - 10 \cdot \langle i \rangle \cdot \#7) \\
& + 10^{-12}(\#4 - \#1 \cdot \#7 - 10 \cdot \langle i \rangle \cdot \#8) + 10^{-16}(-\#1 \cdot \#8),
\end{aligned}$$

where $\langle i \rangle$ stands for the 10^5 digit of Q_A , which is 0 or 1, and $\#1$, $\#2$, *etc.* are the parameters of either auxiliary. The factors of 10 come from the fact that $Q_A = 10 \cdot 10^4 \cdot \langle i \rangle + \#1$. As usual, to combine all the terms, we need to choose some shifts which must ensure that the number of digits of the second, third, and fourth terms are each fixed. Here, the positive contributions are at most 10^8 and the negative contributions can go up to 10^9 . Indeed, for the auxiliary with $\langle i \rangle = 1$, $\#1$ is at most 80000, leading to contributions of at worst $-8 \cdot 10^8 4$, while the other negative term is very small $< 10^6$ (except in the first expression, where we don't care about the number of digits); for the auxiliary with $\langle i \rangle = 0$, $\#1$ can go up to 99999, but there is no other negative term. Hence, a good choice is $2 \cdot 10^9$, which produces totals in the range $[10^9, 2.1 \cdot 10^9]$. We are flirting with $\text{T}_{\text{E}}\text{X}$'s limits once more.

```

19433 \cs_new:Npn \__fp_div_significand_calc:wnnnnnnn 1#1

```

```

19434 {
19435   \if_meaning:w 1 #1
19436   \exp_after:wN \__fp_div_significand_calc_i:wwnnnnnnnn
19437   \else:
19438     \exp_after:wN \__fp_div_significand_calc_ii:wwnnnnnnnn
19439   \fi:
19440 }
19441 \cs_new:Npn \__fp_div_significand_calc_i:wwnnnnnnnn
19442   #1; #2;#3#4 #5#6#7#8 #9
19443   {
19444     1 1 #1
19445     #9 \exp_after:wN ;
19446     \int_value:w \__fp_int_eval:w \c__fp_Bigg_leading_shift_int
19447       + #2 - #1 * #5 - #5#60
19448     \exp_after:wN \__fp_pack_Bigg:NNNNNNw
19449     \int_value:w \__fp_int_eval:w \c__fp_Bigg_middle_shift_int
19450       + #3 - #1 * #6 - #70
19451     \exp_after:wN \__fp_pack_Bigg:NNNNNNw
19452     \int_value:w \__fp_int_eval:w \c__fp_Bigg_middle_shift_int
19453       + #4 - #1 * #7 - #80
19454     \exp_after:wN \__fp_pack_Bigg:NNNNNNw
19455     \int_value:w \__fp_int_eval:w \c__fp_Bigg_trailing_shift_int
19456       - #1 * #8 ;
19457     {#5}{#6}{#7}{#8}
19458   }
19459 \cs_new:Npn \__fp_div_significand_calc_ii:wwnnnnnnnn
19460   #1; #2;#3#4 #5#6#7#8 #9
19461   {
19462     1 0 #1
19463     #9 \exp_after:wN ;
19464     \int_value:w \__fp_int_eval:w \c__fp_Bigg_leading_shift_int
19465       + #2 - #1 * #5
19466     \exp_after:wN \__fp_pack_Bigg:NNNNNNw
19467     \int_value:w \__fp_int_eval:w \c__fp_Bigg_middle_shift_int
19468       + #3 - #1 * #6
19469     \exp_after:wN \__fp_pack_Bigg:NNNNNNw
19470     \int_value:w \__fp_int_eval:w \c__fp_Bigg_middle_shift_int
19471       + #4 - #1 * #7
19472     \exp_after:wN \__fp_pack_Bigg:NNNNNNw
19473     \int_value:w \__fp_int_eval:w \c__fp_Bigg_trailing_shift_int
19474       - #1 * #8 ;
19475     {#5}{#6}{#7}{#8}
19476   }

```

(End definition for __fp_div_significand_calc:wwnnnnnnnn, __fp_div_significand_calc_i:wwnnnnnnnn,
and __fp_div_significand_calc_ii:wwnnnnnnnn.)

__fp_div_significand_ii:wwn $\langle y \rangle$; $\langle B_1 \rangle$; $\{\langle B_2 \rangle\}$ $\{\langle B_3 \rangle\}$ $\{\langle B_4 \rangle\}$ $\{\langle Z_1 \rangle\}$
 $\{\langle Z_2 \rangle\}$ $\{\langle Z_3 \rangle\}$ $\{\langle Z_4 \rangle\}$ $\langle continuations \rangle$ $\langle sign \rangle$

Compute Q_B by evaluating $\langle B_1 \rangle \langle B_2 \rangle 0 / y - 1$. The result is output to the left, in an
 $\backslash_fp_int_eval:w$ which we start now. Once that is evaluated (and the other Q_i also,
since later expansions are triggered by this one), a packing auxiliary takes care of placing
the digits of Q_B in an appropriate way for the final addition to obtain Q . This auxiliary
is also used to compute Q_C and Q_D with the inputs C and D instead of B .

```

19477 \cs_new:Npn \__fp_div_significand_ii:wwn #1; #2;#3
19478 {
19479   \exp_after:wN \__fp_div_significand_pack:NNN
19480   \int_value:w \__fp_int_eval:w
19481   \exp_after:wN \__fp_div_significand_calc:wwnnnnnnn
19482   \int_value:w \__fp_int_eval:w 999999 + #2 #3 0 / #1 ; #2 #3 ;
19483 }

```

(End definition for __fp_div_significand_ii:wwn.)

```

\__fp_div_significand_iii:wwnnnnn <y> ; <E1> ; {<E2>} {<E3>} {<E4>}
{<Z1>} {<Z2>} {<Z3>} {<Z4>} <sign>

```

We compute $P \simeq 2E/Z$ by rounding $2E_1E_2/Z_1Z_2$. Note the first 0, which multiplies Q_D by 10: we later add (roughly) $5 \cdot P$, which amounts to adding $P/2 \simeq E/Z$ to Q_D , the appropriate correction from a hypothetical Q_E .

```

19484 \cs_new:Npn \__fp_div_significand_iii:wwnnnnn #1; #2;#3#4#5 #6#7
19485 {
19486   0
19487   \exp_after:wN \__fp_div_significand_iv:wwnnnnnnn
19488   \int_value:w \__fp_int_eval:w ( 2 * #2 #3 ) / #6 #7 ; % <- P
19489   #2 ; {#3} {#4} {#5}
19490   {#6} {#7}
19491 }

```

(End definition for __fp_div_significand_iii:wwnnnnn.)

```

\__fp_div_significand_iv:wwnnnnnnn <P> ; <E1> ; {<E2>} {<E3>} {<E4>}
\__fp_div_significand_v:NNw {<Z1>} {<Z2>} {<Z3>} {<Z4>} <sign>
\__fp_div_significand_vi:Nw

```

This adds to the current expression $(10^7 + 10 \cdot Q_D)$ a contribution of $5 \cdot P + \text{sign}(T)$ with $T = 2E - PZ$. This amounts to adding $P/2$ to Q_D , with an extra *<rounding>* digit. This *<rounding>* digit is 0 or 5 if T does not contribute, *i.e.*, if $0 = T = 2E - PZ$, in other words if $10^{16}A/Z$ is an integer or half-integer. Otherwise it is in the appropriate range, $[1, 4]$ or $[6, 9]$. This is precise enough for rounding purposes (in any mode).

It seems an overkill to compute T exactly as I do here, but I see no faster way right now.

Once more, we need to be careful and show that the calculation $\#1 \cdot \#6\#7$ below does not cause an overflow: naively, P can be up to 35, and $\#6\#7$ up to 10^8 , but both cannot happen simultaneously. To show that things are fine, we split in two (non-disjoint) cases.

- For $P < 10$, the product obeys $P \cdot \#6\#7 < 10^8 \cdot P < 10^9$.
- For large $P \geq 3$, the rounding error on P , which is at most 1, is less than a factor of 2, hence $P \leq 4E/Z$. Also, $\#6\#7 \leq 10^8 \cdot Z$, hence $P \cdot \#6\#7 \leq 4E \cdot 10^8 < 10^9$.

Both inequalities could be made tighter if needed.

Note however that $P \cdot \#8\#9$ may overflow, since the two factors are now independent, and the result may reach $3.5 \cdot 10^9$. Thus we compute the two lower levels separately. The rest is standard, except that we use $+$ as a separator (ending integer expressions explicitly). T is negative if the first character is $-$, it is positive if the first character is neither 0 nor $-$. It is also positive if the first character is 0 and second argument of $__fp_div_significand_vi:Nw$, a sum of several terms, is also zero. Otherwise, there was an exact agreement: $T = 0$.

```

19492 \cs_new:Npn \__fp_div_significand_iv:wwnnnnnnn #1; #2;#3#4#5 #6#7#8#9
19493 {
19494   + 5 * #1
19495   \exp_after:wN \__fp_div_significand_vi:Nw
19496   \int_value:w \__fp_int_eval:w -20 + 2*#2#3 - #1*#6#7 +
19497   \exp_after:wN \__fp_div_significand_v:NN
19498   \int_value:w \__fp_int_eval:w 199980 + 2*#4 - #1*#8 +
19499   \exp_after:wN \__fp_div_significand_v:NN
19500   \int_value:w \__fp_int_eval:w 200000 + 2*#5 - #1*#9 ;
19501 }
19502 \cs_new:Npn \__fp_div_significand_v:NN #1#2 { #1#2 \__fp_int_eval_end: + }
19503 \cs_new:Npn \__fp_div_significand_vi:Nw #1#2;
19504 {
19505   \if_meaning:w 0 #1
19506     \if_int_compare:w \__fp_int_eval:w #2 > 0 + 1 \fi:
19507   \else:
19508     \if_meaning:w - #1 - \else: + \fi: 1
19509   \fi:
19510   ;
19511 }

```

(End definition for __fp_div_significand_iv:wwnnnnnnn, __fp_div_significand_v:NNw, and __fp_div_significand_vi:Nw.)

__fp_div_significand_pack:NNN

At this stage, we are in the following situation: \TeX is in the process of expanding several integer expressions, thus functions at the bottom expand before those above.

$$\begin{aligned} & _ _ \text{fp_div_significand_test_o:w } 10^6 + Q_A _ _ \text{fp_div_significand_} \\ & \text{pack:NNN } 10^6 + Q_B _ _ \text{fp_div_significand_pack:NNN } 10^6 + Q_C _ _ \text{fp_} \\ & \text{div_significand_pack:NNN } 10^7 + 10 \cdot Q_D + 5 \cdot P + \varepsilon ; \langle \text{sign} \rangle \end{aligned}$$

Here, $\varepsilon = \text{sign}(T)$ is 0 in case $2E = PZ$, 1 in case $2E > PZ$, which means that P was the correct value, but not with an exact quotient, and -1 if $2E < PZ$, *i.e.*, P was an overestimate. The packing function we define now does nothing special: it removes the 10^6 and carries two digits (for the 10^5 's and the 10^4 's).

```

19512 \cs_new:Npn \__fp_div_significand_pack:NNN 1 #1 #2 { + #1 #2 ; }

```

(End definition for __fp_div_significand_pack:NNN.)

__fp_div_significand_test_o:w

__fp_div_significand_test_o:w 1 0 $\langle 5d \rangle$; $\langle 4d \rangle$; $\langle 4d \rangle$; $\langle 5d \rangle$; $\langle \text{sign} \rangle$

The reason we know that the first two digits are 1 and 0 is that the final result is known to be between 0.1 (inclusive) and 10, hence \widetilde{Q}_A (the tilde denoting the contribution from the other Q_i) is at most 99999, and $10^6 + \widetilde{Q}_A = 10 \dots$.

It is now time to round. This depends on how many digits the final result will have.

```

19513 \cs_new:Npn \__fp_div_significand_test_o:w 10 #1
19514 {
19515   \if_meaning:w 0 #1
19516     \exp_after:wN \__fp_div_significand_small_o:wwwNNNNwN
19517   \else:
19518     \exp_after:wN \__fp_div_significand_large_o:wwwNNNNwN
19519   \fi:
19520   #1
19521 }

```

(End definition for __fp_div_significand_test_o:w.)

```

\__fp_div_significand_small_o:wwwNNNNwN \__fp_div_significand_small_o:wwwNNNNwN 0 <4d> ; <4d> ; <4d> ; <5d>
; <final sign>

```

Standard use of the functions `__fp_basics_pack_low:NNNNw` and `__fp_basics_pack_high:NNNNw`. We finally get to use the *<final sign>* which has been sitting there for a while.

```

19522 \cs_new:Npn \__fp_div_significand_small_o:wwwNNNNwN
19523 0 #1; #2; #3; #4#5#6#7#8; #9
19524 {
19525   \exp_after:wN \__fp_basics_pack_high:NNNNw
19526   \int_value:w \__fp_int_eval:w 1 #1#2
19527   \exp_after:wN \__fp_basics_pack_low:NNNNw
19528   \int_value:w \__fp_int_eval:w 1 #3#4#5#6#7
19529   + \__fp_round:NNN #9 #7 #8
19530   \exp_after:wN ;
19531 }

```

(End definition for `__fp_div_significand_small_o:wwwNNNNwN`.)

```

\__fp_div_significand_large_o:wwwNNNNwN \__fp_div_significand_large_o:wwwNNNNwN <5d> ; <4d> ; <4d> ; <5d> ;
<sign>

```

We know that the final result cannot reach 10, hence `1#1#2`, together with contributions from the level below, cannot reach $2 \cdot 10^9$. For rounding, we build the *<rounding digit>* from the last two of our 18 digits.

```

19532 \cs_new:Npn \__fp_div_significand_large_o:wwwNNNNwN
19533 #1; #2; #3; #4#5#6#7#8; #9
19534 {
19535   + 1
19536   \exp_after:wN \__fp_basics_pack_weird_high:NNNNNNw
19537   \int_value:w \__fp_int_eval:w 1 #1 #2
19538   \exp_after:wN \__fp_basics_pack_weird_low:NNNNw
19539   \int_value:w \__fp_int_eval:w 1 #3 #4 #5 #6 +
19540   \exp_after:wN \__fp_round:NNN
19541   \exp_after:wN #9
19542   \exp_after:wN #6
19543   \int_value:w \__fp_round_digit:Nw #7 #8 ;
19544   \exp_after:wN ;
19545 }

```

(End definition for `__fp_div_significand_large_o:wwwNNNNwN`.)

31.4 Square root

`__fp_sqrt_o:w` Zeros are unchanged: $\sqrt{-0} = -0$ and $\sqrt{+0} = +0$. Negative numbers (other than -0) have no real square root. Positive infinity, and `nan`, are unchanged. Finally, for normal positive numbers, there is some work to do.

```

19546 \cs_new:Npn \__fp_sqrt_o:w #1 \s__fp \__fp_chk:w #2#3#4; @
19547 {
19548   \if_meaning:w 0 #2 \__fp_case_return_same_o:w \fi:
19549   \if_meaning:w 2 #3
19550     \__fp_case_use:nw { \__fp_invalid_operation_o:nw { sqrt } }
19551   \fi:
19552   \if_meaning:w 1 #2 \else: \__fp_case_return_same_o:w \fi:
19553   \__fp_sqrt_npos_o:w

```

```

19554   \s__fp \__fp_chk:w #2 #3 #4;
19555   }

```

(End definition for __fp_sqrt_o:w.)

```

\__fp_sqrt_npos_o:w
\__fp_sqrt_npos_auxi_o:wNnnN
\__fp_sqrt_npos_auxii_o:wNNNNNNNN

```

Prepare __fp_sanitize:Nw to receive the final sign 0 (the result is always positive) and the exponent, equal to half of the exponent #1 of the argument. If the exponent #1 is even, find a first approximation of the square root of the significand $10^8 a_1 + a_2 = 10^8 \#2\#3 + \#4\#5$ through Newton's method, starting at $x = 57234133 \simeq 10^{7.75}$. Otherwise, first shift the significand of the argument by one digit, getting $a'_1 \in [10^6, 10^7)$ instead of $[10^7, 10^8)$, then use Newton's method starting at $17782794 \simeq 10^{7.25}$.

```

19556 \cs_new:Npn \__fp_sqrt_npos_o:w \s__fp \__fp_chk:w 1 0 #1#2#3#4#5;
19557 {
19558   \exp_after:wN \__fp_sanitize:Nw
19559   \exp_after:wN 0
19560   \int_value:w \__fp_int_eval:w
19561   \if_int_odd:w #1 \exp_stop_f:
19562     \exp_after:wN \__fp_sqrt_npos_auxi_o:wNnnN
19563     \fi:
19564     #1 / 2
19565     \__fp_sqrt_Newton_o:wN 56234133; 0; {#2#3} {#4#5} 0
19566   }
19567 \cs_new:Npn \__fp_sqrt_npos_auxi_o:wNnnN #1 / 2 #2; 0; #3#4#5
19568 {
19569   ( #1 + 1 ) / 2
19570   \__fp_pack_eight:wNNNNNNNN
19571   \__fp_sqrt_npos_auxii_o:wNNNNNNNN
19572   ;
19573   0 #3 #4
19574 }
19575 \cs_new:Npn \__fp_sqrt_npos_auxii_o:wNNNNNNNN #1; #2#3#4#5#6#7#8#9
19576 { \__fp_sqrt_Newton_o:wN 17782794; 0; {#1} {#2#3#4#5#6#7#8#9} }

```

(End definition for __fp_sqrt_npos_o:w, __fp_sqrt_npos_auxi_o:wNnnN, and __fp_sqrt_npos_auxii_o:wNNNNNNNN.)

```

\__fp_sqrt_Newton_o:wN

```

Newton's method maps $x \mapsto [(x + [10^8 a_1/x])/2]$ in each iteration, where $[b/c]$ denotes ε -TeX's division. This division rounds the real number b/c to the closest integer, rounding ties away from zero, hence when c is even, $b/c - 1/2 + 1/c \leq [b/c] \leq b/c + 1/2$ and when c is odd, $b/c - 1/2 + 1/(2c) \leq [b/c] \leq b/c + 1/2 - 1/(2c)$. For all c , $b/c - 1/2 + 1/(2c) \leq [b/c] \leq b/c + 1/2$.

Let us prove that the method converges when implemented with ε -TeX integer division, for any $10^6 \leq a_1 < 10^8$ and starting value $10^6 \leq x < 10^8$. Using the inequalities above and the arithmetic-geometric inequality $(x + t)/2 \geq \sqrt{xt}$ for $t = 10^8 a_1/x$, we find

$$x' = \left\lceil \frac{x + [10^8 a_1/x]}{2} \right\rceil \geq \frac{x + 10^8 a_1/x - 1/2 + 1/(2x)}{2} \geq \sqrt{10^8 a_1} - \frac{1}{4} + \frac{1}{4x}.$$

After any step of iteration, we thus have $\delta = x - \sqrt{10^8 a_1} \geq -0.25 + 0.25 \cdot 10^{-8}$. The new difference $\delta' = x' - \sqrt{10^8 a_1}$ after one step is bounded above as

$$x' - \sqrt{10^8 a_1} \leq \frac{x + 10^8 a_1/x + 1/2}{2} + \frac{1}{2} - \sqrt{10^8 a_1} \leq \frac{\delta}{2} \frac{\delta}{\sqrt{10^8 a_1} + \delta} + \frac{3}{4}.$$

For $\delta > 3/2$, this last expression is $\leq \delta/2 + 3/4 < \delta$, hence δ decreases at each step: since all x are integers, δ must reach a value $-1/4 < \delta \leq 3/2$. In this range of values, we get $\delta' \leq \frac{3}{4} \frac{3}{2\sqrt{10^8 a_1}} + \frac{3}{4} \leq 0.75 + 1.125 \cdot 10^{-7}$. We deduce that the difference $\delta = x - \sqrt{10^8 a_1}$ eventually reaches a value in the interval $[-0.25 + 0.25 \cdot 10^{-8}, 0.75 + 11.25 \cdot 10^{-8}]$, whose width is $1 + 11 \cdot 10^{-8}$. The corresponding interval for x may contain two integers, hence x might oscillate between those two values.

However, the fact that $x \mapsto x - 1$ and $x - 1 \mapsto x$ puts stronger constraints, which are not compatible: the first implies

$$x + [10^8 a_1 / x] \leq 2x - 2$$

hence $10^8 a_1 / x \leq x - 3/2$, while the second implies

$$x - 1 + [10^8 a_1 / (x - 1)] \geq 2x - 1$$

hence $10^8 a_1 / (x - 1) \geq x - 1/2$. Combining the two inequalities yields $x^2 - 3x/2 \geq 10^8 a_1 \geq x - 3x/2 + 1/2$, which cannot hold. Therefore, the iteration always converges to a single integer x . To stop the iteration when two consecutive results are equal, the function `_fp_sqrt_Newton_o:wnn` receives the newly computed result as #1, the previous result as #2, and a_1 as #3. Note that ε -TeX combines the computation of a multiplication and a following division, thus avoiding overflow in `#3 * 100000000 / #1`. In any case, the result is within $[10^7, 10^8]$.

```

19577 \cs_new:Npn \_fp_sqrt_Newton_o:wnn #1; #2; #3
19578 {
19579   \if_int_compare:w #1 = #2 \exp_stop_f:
19580     \exp_after:wN \_fp_sqrt_auxi_o:NNNNwnnnN
19581     \int_value:w \_fp_int_eval:w 9999 9999 +
19582     \exp_after:wN \_fp_use_none_until_s:w
19583   \fi:
19584   \exp_after:wN \_fp_sqrt_Newton_o:wnn
19585   \int_value:w \_fp_int_eval:w (#1 + #3 * 1 0000 0000 / #1) / 2 ;
19586   #1; {#3}
19587 }

```

(End definition for `_fp_sqrt_Newton_o:wnn`.)

`_fp_sqrt_auxi_o:NNNNwnnnN` This function is followed by $10^8 + x - 1$, which has 9 digits starting with 1, then ; $\{a_1\} \{a_2\} \{a'\}$. Here, $x \simeq \sqrt{10^8 a_1}$ and we want to estimate the square root of $a = 10^{-8} a_1 + 10^{-16} a_2 + 10^{-17} a'$. We set up an initial underestimate

$$y = (x - 1)10^{-8} + 0.2499998875 \cdot 10^{-8} \lesssim \sqrt{a}.$$

From the inequalities shown earlier, we know that $y \leq \sqrt{10^{-8} a_1} \leq \sqrt{a}$ and that $\sqrt{10^{-8} a_1} \leq y + 10^{-8} + 11 \cdot 10^{-16}$ hence (using $0.1 \leq y \leq \sqrt{a} \leq 1$)

$$a - y^2 \leq 10^{-8} a_1 + 10^{-8} - y^2 \leq (y + 10^{-8} + 11 \cdot 10^{-16})^2 - y^2 + 10^{-8} < 3.2 \cdot 10^{-8},$$

and $\sqrt{a} - y = (a - y^2) / (\sqrt{a} + y) \leq 16 \cdot 10^{-8}$. Next, `_fp_sqrt_auxii_o:NNnnnnnnN` is called several times to get closer and closer underestimates of \sqrt{a} . By construction, the underestimates y are always increasing, $a - y^2 < 3.2 \cdot 10^{-8}$ for all. Also, $y < 1$.

```

19588 \cs_new:Npn \_fp_sqrt_auxi_o:NNNNwnnnN 1 #1#2#3#4#5;
19589 {

```

```

19590 \_fp_sqrt_auxii_o:NnnnnnnnN
19591 \_fp_sqrt_auxiii_o:wnnnnnnnn
19592 {#1#2#3#4} {#5} {2499} {9988} {7500}
19593 }

```

(End definition for _fp_sqrt_auxi_o:NNNNwnnnN.)

_fp_sqrt_auxii_o:NnnnnnnnN

This receives a continuation function #1, then five blocks of 4 digits for y , then two 8-digit blocks and a single digit for a . A common estimate of $\sqrt{a} - y = (a - y^2)/(\sqrt{a} + y)$ is $(a - y^2)/(2y)$, which leads to alternating overestimates and underestimates. We tweak this, to only work with underestimates (no need then to worry about signs in the computation). Each step finds the largest integer $j \leq 6$ such that $10^{4j}(a - y^2) < 2 \cdot 10^8$, then computes the integer (with ε -TeX's rounding division)

$$10^{4j}z = \left[\left(\lfloor 10^{4j}(a - y^2) \rfloor - 257 \right) \cdot (0.5 \cdot 10^8) \right] / \lfloor 10^8 y + 1 \rfloor.$$

The choice of j ensures that $10^{4j}z < 2 \cdot 10^8 \cdot 0.5 \cdot 10^8 / 10^7 = 10^9$, thus $10^9 + 10^{4j}z$ has exactly 10 digits, does not overflow TeX's integer range, and starts with 1. Incidentally, since all $a - y^2 \leq 3.2 \cdot 10^{-8}$, we know that $j \geq 3$.

Let us show that z is an underestimate of $\sqrt{a} - y$. On the one hand, $\sqrt{a} - y \leq 16 \cdot 10^{-8}$ because this holds for the initial y and values of y can only increase. On the other hand, the choice of j implies that $\sqrt{a} - y \leq 5(\sqrt{a} + y)(\sqrt{a} - y) = 5(a - y^2) < 10^{9-4j}$. For $j = 3$, the first bound is better, while for larger j , the second bound is better. For all $j \in [3, 6]$, we find $\sqrt{a} - y < 16 \cdot 10^{-2j}$. From this, we deduce that

$$10^{4j}(\sqrt{a} - y) = \frac{10^{4j}(a - y^2 - (\sqrt{a} - y)^2)}{2y} \geq \frac{\lfloor 10^{4j}(a - y^2) \rfloor - 257}{2 \cdot 10^{-8} \lfloor 10^8 y + 1 \rfloor} + \frac{1}{2}$$

where we have replaced the bound $10^{4j}(16 \cdot 10^{-2j}) = 256$ by 257 and extracted the corresponding term $1/(2 \cdot 10^{-8} \lfloor 10^8 y + 1 \rfloor) \geq 1/2$. Given that ε -TeX's integer division obeys $\lfloor b/c \rfloor \leq b/c + 1/2$, we deduce that $10^{4j}z \leq 10^{4j}(\sqrt{a} - y)$, hence $y + z \leq \sqrt{a}$ is an underestimate of \sqrt{a} , as claimed. One implementation detail: because the computation involves $-4*4 - 2*3*5 - 2*2*6$ which may be as low as $-5 \cdot 10^8$, we need to use the `pack_big` functions, and the big shifts.

```

19594 \cs_new:Npn \_fp_sqrt_auxii_o:NnnnnnnnN #1 #2#3#4#5#6 #7#8#9
19595 {
19596   \exp_after:wN #1
19597   \int_value:w \_fp_int_eval:w \c\_fp_big_leading_shift_int
19598   + #7 - #2 * #2
19599   \exp_after:wN \_fp_pack_big:NNNNNNw
19600   \int_value:w \_fp_int_eval:w \c\_fp_big_middle_shift_int
19601   - 2 * #2 * #3
19602   \exp_after:wN \_fp_pack_big:NNNNNNw
19603   \int_value:w \_fp_int_eval:w \c\_fp_big_middle_shift_int
19604   + #8 - #3 * #3 - 2 * #2 * #4
19605   \exp_after:wN \_fp_pack_big:NNNNNNw
19606   \int_value:w \_fp_int_eval:w \c\_fp_big_middle_shift_int
19607   - 2 * #3 * #4 - 2 * #2 * #5
19608   \exp_after:wN \_fp_pack_big:NNNNNNw
19609   \int_value:w \_fp_int_eval:w \c\_fp_big_middle_shift_int
19610   + #9 000 0000 - #4 * #4 - 2 * #3 * #5 - 2 * #2 * #6
19611   \exp_after:wN \_fp_pack_big:NNNNNNw

```



```

19612      \int_value:w \_fp_int_eval:w \c\_fp\_big\_middle\_shift\_int
19613      - 2 * #4 * #5 - 2 * #3 * #6
19614      \exp_after:wN \_fp\_pack\_big:NNNNNNw
19615      \int_value:w \_fp_int_eval:w \c\_fp\_big\_middle\_shift\_int
19616      - #5 * #5 - 2 * #4 * #6
19617      \exp_after:wN \_fp\_pack\_big:NNNNNNw
19618      \int_value:w \_fp_int_eval:w
19619      \c\_fp\_big\_middle\_shift\_int
19620      - 2 * #5 * #6
19621      \exp_after:wN \_fp\_pack\_big:NNNNNNw
19622      \int_value:w \_fp_int_eval:w
19623      \c\_fp\_big\_trailing\_shift\_int
19624      - #6 * #6 ;
19625      % (
19626      - 257 ) * 5000 0000 / (#2#3 + 1) + 10 0000 0000 ;
19627      {#2}{#3}{#4}{#5}{#6} {#7}{#8}#9
19628      }

```

(End definition for _fp_sqrt_auxii_o:NnnnnnnnnN.)

```

\_fp\_sqrt\_auxiii\_o:wnnnnnnnnn
\_fp\_sqrt\_auxiv\_o:NNNNNNw
\_fp\_sqrt\_auxv\_o:NNNNNNw
\_fp\_sqrt\_auxvi\_o:NNNNNNw
\_fp\_sqrt\_auxvii\_o:NNNNNNw

```

We receive here the difference $a - y^2 = d = \sum_i d_i \cdot 10^{-4i}$, as $\langle d_2 \rangle$; $\{\langle d_3 \rangle\} \dots \{\langle d_{10} \rangle\}$, where each block has 4 digits, except $\langle d_2 \rangle$. This function finds the largest $j \leq 6$ such that $10^{4j}(a - y^2) < 2 \cdot 10^8$, then leaves an open parenthesis and the integer $\lfloor 10^{4j}(a - y^2) \rfloor$ in an integer expression. The closing parenthesis is provided by the caller _fp_sqrt_auxii_o:NnnnnnnnnN, which completes the expression

$$10^{4j}z = \left[(\lfloor 10^{4j}(a - y^2) \rfloor - 257) \cdot (0.5 \cdot 10^8) / \lfloor 10^8 y + 1 \rfloor \right]$$

for an estimate of $10^{4j}(\sqrt{a} - y)$. If $d_2 \geq 2$, $j = 3$ and the **auxiv** auxiliary receives $10^{12}z$. If $d_2 \leq 1$ but $10^4 d_2 + d_3 \geq 2$, $j = 4$ and the **auxv** auxiliary is called, and receives $10^{16}z$, and so on. In all those cases, the **auxviii** auxiliary is set up to add z to y , then go back to the **auxii** step with continuation **auxiii** (the function we are currently describing). The maximum value of j is 6, regardless of whether $10^{12}d_2 + 10^8 d_3 + 10^4 d_4 + d_5 \geq 1$. In this last case, we detect when $10^{24}z < 10^7$, which essentially means $\sqrt{a} - y \lesssim 10^{-17}$: once this threshold is reached, there is enough information to find the correctly rounded \sqrt{a} with only one more call to _fp_sqrt_auxii_o:NnnnnnnnnN. Note that the iteration cannot be stuck before reaching $j = 6$, because for $j < 6$, one has $2 \cdot 10^8 \leq 10^{4(j+1)}(a - y^2)$, hence

$$10^{4j}z \geq \frac{(20000 - 257)(0.5 \cdot 10^8)}{\lfloor 10^8 y + 1 \rfloor} \geq (20000 - 257) \cdot 0.5 > 0.$$

```

19629 \cs_new:Npn \_fp\_sqrt\_auxiii\_o:wnnnnnnnnn
19630      #1; #2#3#4#5#6#7#8#9
19631      {
19632      \if_int_compare:w #1 > 1 \exp_stop_f:
19633      \exp_after:wN \_fp\_sqrt\_auxiv\_o:NNNNNNw
19634      \int_value:w \_fp_int_eval:w (#1#2 %)
19635      \else:
19636      \if_int_compare:w #1#2 > 1 \exp_stop_f:
19637      \exp_after:wN \_fp\_sqrt\_auxv\_o:NNNNNNw
19638      \int_value:w \_fp_int_eval:w (#1#2#3 %)
19639      \else:
19640      \if_int_compare:w #1#2#3 > 1 \exp_stop_f:
19641      \exp_after:wN \_fp\_sqrt\_auxvi\_o:NNNNNNw

```

```

19642         \int_value:w \__fp_int_eval:w (#1#2#3#4 %)
19643     \else:
19644         \exp_after:wN \__fp_sqrt_auxvii_o:NNNNNw
19645         \int_value:w \__fp_int_eval:w (#1#2#3#4#5 %)
19646     \fi:
19647 \fi:
19648 \fi:
19649 }
19650 \cs_new:Npn \__fp_sqrt_auxiv_o:NNNNNw #1#2#3#4#5#6;
19651 { \__fp_sqrt_auxviii_o:nnnnnnnn {#1#2#3#4#5#6} {00000000} }
19652 \cs_new:Npn \__fp_sqrt_auxv_o:NNNNNw #1#2#3#4#5#6;
19653 { \__fp_sqrt_auxviii_o:nnnnnnnn {000#1#2#3#4#5} {#60000} }
19654 \cs_new:Npn \__fp_sqrt_auxvi_o:NNNNNw #1#2#3#4#5#6;
19655 { \__fp_sqrt_auxviii_o:nnnnnnnn {0000000#1} {#2#3#4#5#6} }
19656 \cs_new:Npn \__fp_sqrt_auxvii_o:NNNNNw #1#2#3#4#5#6;
19657 {
19658     \if_int_compare:w #1#2 = 0 \exp_stop_f:
19659     \exp_after:wN \__fp_sqrt_auxx_o:Nnnnnnnnn
19660     \fi:
19661     \__fp_sqrt_auxviii_o:nnnnnnnn {00000000} {000#1#2#3#4#5}
19662 }

```

(End definition for `__fp_sqrt_auxiii_o:wnnnnnnnn` and others.)

`__fp_sqrt_auxviii_o:nnnnnnnn` Simply add the two 8-digit blocks of z , aligned to the last four of the five 4-digit blocks of y , then call the `auxii` auxiliary to evaluate $y'^2 = (y + z)^2$.

```

19663 \cs_new:Npn \__fp_sqrt_auxviii_o:nnnnnnnn #1#2 #3#4#5#6#7
19664 {
19665     \exp_after:wN \__fp_sqrt_auxix_o:wnwnnw
19666     \int_value:w \__fp_int_eval:w #3
19667     \exp_after:wN \__fp_basics_pack_low:NNNNNw
19668     \int_value:w \__fp_int_eval:w #1 + 1#4#5
19669     \exp_after:wN \__fp_basics_pack_low:NNNNNw
19670     \int_value:w \__fp_int_eval:w #2 + 1#6#7 ;
19671 }
19672 \cs_new:Npn \__fp_sqrt_auxix_o:wnwnnw #1; #2#3; #4#5;
19673 {
19674     \__fp_sqrt_auxii_o:NnnnnnnnnN
19675     \__fp_sqrt_auxiii_o:wnnnnnnnnn {#1}{#2}{#3}{#4}{#5}
19676 }

```

(End definition for `__fp_sqrt_auxviii_o:nnnnnnnn` and `__fp_sqrt_auxix_o:wnwnnw`.)

`__fp_sqrt_auxx_o:Nnnnnnnnn` At this stage, $j = 6$ and $10^{24}z < 10^7$, hence
`__fp_sqrt_auxxi_o:wnnnN`

$$10^7 + 1/2 > 10^{24}z + 1/2 \geq (10^{24}(a - y^2) - 258) \cdot (0.5 \cdot 10^8) / (10^8y + 1),$$

then $10^{24}(a - y^2) - 258 < 2(10^7 + 1/2)(y + 10^{-8})$, and

$$10^{24}(a - y^2) < (10^7 + 1290.5)(1 + 10^{-8}/y)(2y) < (10^7 + 1290.5)(1 + 10^{-7})(y + \sqrt{a}),$$

which finally implies $0 \leq \sqrt{a} - y < 0.2 \cdot 10^{-16}$. In particular, y is an underestimate of \sqrt{a} and $y + 0.5 \cdot 10^{-16}$ is a (strict) overestimate. There is at exactly one multiple m of $0.5 \cdot 10^{-16}$ in the interval $[y, y + 0.5 \cdot 10^{-16})$. If $m^2 > a$, then the square root is inexact and

is obtained by rounding $m - \epsilon$ to a multiple of 10^{-16} (the precise shift $0 < \epsilon < 0.5 \cdot 10^{-16}$ is irrelevant for rounding). If $m^2 = a$ then the square root is exactly m , and there is no rounding. If $m^2 < a$ then we round $m + \epsilon$. For now, discard a few irrelevant arguments #1, #2, #3, and find the multiple of $0.5 \cdot 10^{-16}$ within $[y, y + 0.5 \cdot 10^{-16})$; rather, only the last 4 digits #8 of y are considered, and we do not perform any carry yet. The `auxxi` auxiliary sets up `auxii` with a continuation function `auxxii` instead of `auxiii` as before. To prevent `auxii` from giving a negative results $a - m^2$, we compute $a + 10^{-16} - m^2$ instead, always positive since $m < \sqrt{a} + 0.5 \cdot 10^{-16}$ and $a \leq 1 - 10^{-16}$.

```

19677 \cs_new:Npn \__fp_sqrt_auxx_o:Nnnnnnnn #1#2#3 #4#5#6#7#8
19678 {
19679   \exp_after:wN \__fp_sqrt_auxxi_o:wwnnN
19680   \int_value:w \__fp_int_eval:w
19681     (#8 + 2499) / 5000 * 5000 ;
19682   {#4} {#5} {#6} {#7} ;
19683 }
19684 \cs_new:Npn \__fp_sqrt_auxxi_o:wwnnN #1; #2; #3#4#5
19685 {
19686   \__fp_sqrt_auxii_o:NnnnnnnnN
19687   \__fp_sqrt_auxxii_o:nnnnnnnnw
19688   #2 {#1}
19689   {#3} { #4 + 1 } #5
19690 }

```

(End definition for `__fp_sqrt_auxx_o:Nnnnnnnn` and `__fp_sqrt_auxxi_o:wwnnN`.)

`__fp_sqrt_auxxii_o:nnnnnnnnw`
`__fp_sqrt_auxxiii_o:w`

The difference $0 \leq a + 10^{-16} - m^2 \leq 10^{-16} + (\sqrt{a} - m)(\sqrt{a} + m) \leq 2 \cdot 10^{-16}$ was just computed: its first 8 digits vanish, as do the next four, #1, and most of the following four, #2. The guess m is an overestimate if $a + 10^{-16} - m^2 < 10^{-16}$, that is, #1#2 vanishes. Otherwise it is an underestimate, unless $a + 10^{-16} - m^2 = 10^{-16}$ exactly. For an underestimate, call the `auxxiv` function with argument 9998. For an exact result call it with 9999, and for an overestimate call it with 10000.

```

19691 \cs_new:Npn \__fp_sqrt_auxxii_o:nnnnnnnnw 0; #1#2#3#4#5#6#7#8 #9;
19692 {
19693   \if_int_compare:w #1#2 > 0 \exp_stop_f:
19694   \if_int_compare:w #1#2 = 1 \exp_stop_f:
19695   \if_int_compare:w #3#4 = 0 \exp_stop_f:
19696   \if_int_compare:w #5#6 = 0 \exp_stop_f:
19697   \if_int_compare:w #7#8 = 0 \exp_stop_f:
19698     \__fp_sqrt_auxxiii_o:w
19699   \fi:
19700   \fi:
19701   \fi:
19702   \fi:
19703   \exp_after:wN \__fp_sqrt_auxxiv_o:wnnnnnnnN
19704   \int_value:w 9998
19705   \else:
19706     \exp_after:wN \__fp_sqrt_auxxiv_o:wnnnnnnnN
19707     \int_value:w 10000
19708   \fi:
19709   ;
19710 }
19711 \cs_new:Npn \__fp_sqrt_auxxiii_o:w \fi: \fi: \fi: \fi: #1 \fi: ;
19712 {

```

```

19713     \fi: \fi: \fi: \fi: \fi:
19714     \__fp_sqrt_auxxiv_o:wnnnnnnnnN 9999 ;
19715 }

```

(End definition for `__fp_sqrt_auxxii_o:nnnnnnnnw` and `__fp_sqrt_auxxiii_o:w`.)

`__fp_sqrt_auxxiv_o:wnnnnnnnnN`

This receives 9998, 9999 or 10000 as #1 when m is an underestimate, exact, or an overestimate, respectively. Then comes m as five blocks of 4 digits, but where the last block #6 may be 0, 5000, or 10000. In the latter case, we need to add a carry, unless m is an overestimate (#1 is then 10000). Then comes a as three arguments. Rounding is done by `__fp_round:NNN`, whose first argument is the final sign 0 (square roots are positive). We fake its second argument. It should be the last digit kept, but this is only used when ties are “rounded to even”, and only when the result is exactly half-way between two representable numbers rational square roots of numbers with 16 significant digits have: this situation never arises for the square root, as any exact square root of a 16 digit number has at most 8 significant digits. Finally, the last argument is the next digit, possibly shifted by 1 when there are further nonzero digits. This is achieved by `__fp_round_digit:Nw`, which receives (after removal of the 10000’s digit) one of 0000, 0001, 4999, 5000, 5001, or 9999, which it converts to 0, 1, 4, 5, 6, and 9, respectively.

```

19716 \cs_new:Npn \__fp_sqrt_auxxiv_o:wnnnnnnnnN #1; #2#3#4#5#6 #7#8#9
19717 {
19718     \exp_after:wN \__fp_basics_pack_high:NNNNNw
19719     \int_value:w \__fp_int_eval:w 1 0000 0000 + #2#3
19720     \exp_after:wN \__fp_basics_pack_low:NNNNNw
19721     \int_value:w \__fp_int_eval:w 1 0000 0000
19722     + #4#5
19723     \if_int_compare:w #6 > #1 \exp_stop_f: + 1 \fi:
19724     + \exp_after:wN \__fp_round:NNN
19725     \exp_after:wN 0
19726     \exp_after:wN 0
19727     \int_value:w
19728     \exp_after:wN \use_i:nn
19729     \exp_after:wN \__fp_round_digit:Nw
19730     \int_value:w \__fp_int_eval:w #6 + 19999 - #1 ;
19731     \exp_after:wN ;
19732 }

```

(End definition for `__fp_sqrt_auxxiv_o:wnnnnnnnnN`.)

31.5 About the sign and exponent

`__fp_logb_o:w`
`__fp_logb_aux_o:w`

The exponent of a normal number is its *exponent* minus one.

```

19733 \cs_new:Npn \__fp_logb_o:w ? \s__fp \__fp_chk:w #1#2; @
19734 {
19735     \if_case:w #1 \exp_stop_f:
19736         \__fp_case_use:nw
19737         { \__fp_division_by_zero_o:Nnw \c_minus_inf_fp { logb } }
19738     \or: \exp_after:wN \__fp_logb_aux_o:w
19739     \or: \__fp_case_return_o:Nw \c_inf_fp
19740     \else: \__fp_case_return_same_o:w
19741     \fi:
19742     \s__fp \__fp_chk:w #1 #2;
19743 }

```

```

19744 \cs_new:Npn \__fp_logb_aux_o:w \s__fp \__fp_chk:w #1 #2 #3 #4 ;
19745 {
19746   \exp_after:wN \__fp_parse:n \exp_after:wN
19747   { \int_value:w \int_eval:w #3 - 1 \exp_after:wN }
19748 }

```

(End definition for __fp_logb_o:w and __fp_logb_aux_o:w.)

```

\__fp_sign_o:w Find the sign of the floating point: nan, +0, -0, +1 or -1.
\__fp_sign_aux_o:w
19749 \cs_new:Npn \__fp_sign_o:w ? \s__fp \__fp_chk:w #1#2; @
19750 {
19751   \if_case:w #1 \exp_stop_f:
19752     \__fp_case_return_same_o:w
19753   \or: \exp_after:wN \__fp_sign_aux_o:w
19754   \or: \exp_after:wN \__fp_sign_aux_o:w
19755   \else: \__fp_case_return_same_o:w
19756   \fi:
19757   \s__fp \__fp_chk:w #1 #2;
19758 }
19759 \cs_new:Npn \__fp_sign_aux_o:w \s__fp \__fp_chk:w #1 #2 #3 ;
19760 { \exp_after:wN \__fp_set_sign_o:w \exp_after:wN #2 \c_one_fp @ }

```

(End definition for __fp_sign_o:w and __fp_sign_aux_o:w.)

__fp_set_sign_o:w This function is used for the unary minus and for `abs`. It leaves the sign of `nan` invariant, turns negative numbers (sign 2) to positive numbers (sign 0) and positive numbers (sign 0) to positive or negative numbers depending on #1. It also expands after itself in the input stream, just like `__fp_+_o:ww`.

```

19761 \cs_new:Npn \__fp_set_sign_o:w #1 \s__fp \__fp_chk:w #2#3#4; @
19762 {
19763   \exp_after:wN \__fp_exp_after_o:w
19764   \exp_after:wN \s__fp
19765   \exp_after:wN \__fp_chk:w
19766   \exp_after:wN #2
19767   \int_value:w
19768   \if_case:w #3 \exp_stop_f: #1 \or: 1 \or: 0 \fi: \exp_stop_f:
19769   #4;
19770 }

```

(End definition for __fp_set_sign_o:w.)

31.6 Operations on tuples

__fp_tuple_set_sign_o:w Two cases: `abs(<tuple>)` for which #1 is 0 (invalid for tuples) and `-<tuple>` for which #1 is 2. In that case, map over all items in the tuple an auxiliary that dispatches to the type-appropriate sign-flipping function.

```

\__fp_tuple_set_sign_aux_o:Nnw
\__fp_tuple_set_sign_aux_o:w
19771 \cs_new:Npn \__fp_tuple_set_sign_o:w #1
19772 {
19773   \if_meaning:w 2 #1
19774     \exp_after:wN \__fp_tuple_set_sign_aux_o:Nnw
19775   \fi:
19776   \__fp_invalid_operation_o:nw { abs }
19777 }
19778 \cs_new:Npn \__fp_tuple_set_sign_aux_o:Nnw #1#2#3 @

```

```

19779 { \_fp_tuple_map_o:nw \_fp_tuple_set_sign_aux_o:w #3 }
19780 \cs_new:Npn \_fp_tuple_set_sign_aux_o:w #1#2 ;
19781 {
19782   \_fp_change_func_type:NNN #1 \_fp_set_sign_o:w
19783   \_fp_parse_apply_unary_error:NNw
19784   2 #1 #2 ; @
19785 }

```

(End definition for `_fp_tuple_set_sign_o:w`, `_fp_tuple_set_sign_aux_o:Nnw`, and `_fp_tuple_set_sign_aux_o:w`.)

`_fp*_tuple_o:ww` For $\langle number \rangle * \langle tuple \rangle$ and $\langle tuple \rangle * \langle number \rangle$ and $\langle tuple \rangle / \langle number \rangle$, loop through the `_fp_tuple*_o:ww` $\langle tuple \rangle$ some code that multiplies or divides by the appropriate $\langle number \rangle$. Importantly `_fp_tuple/_o:ww` we need to dispatch according to the type, and we make sure to apply the operator in the correct order.

```

19786 \cs_new:cpn { \_fp*_tuple_o:ww } #1 ;
19787 { \_fp_tuple_map_o:nw { \_fp_binary_type_o:Nww * #1 ; } }
19788 \cs_new:cpn { \_fp_tuple*_o:ww } #1 ; #2 ;
19789 { \_fp_tuple_map_o:nw { \_fp_binary_rev_type_o:Nww * #2 ; } #1 ; }
19790 \cs_new:cpn { \_fp_tuple/_o:ww } #1 ; #2 ;
19791 { \_fp_tuple_map_o:nw { \_fp_binary_rev_type_o:Nww / #2 ; } #1 ; }

```

(End definition for `_fp*_tuple_o:ww`, `_fp_tuple*_o:ww`, and `_fp_tuple/_o:ww`.)

`_fp_tuple+_tuple_o:ww` Check the two tuples have the same number of items and map through these a helper `_fp_tuple-_tuple_o:ww` that dispatches appropriately depending on the types. This means $(1,2) + ((1,1),2)$ gives $(\text{nan},4)$.

```

19792 \cs_set_protected:Npn \_fp_tmp:w #1
19793 {
19794   \cs_new:cpn { \_fp_tuple_#1_tuple_o:ww }
19795   \s_fp_tuple \_fp_tuple_chk:w ##1 ;
19796   \s_fp_tuple \_fp_tuple_chk:w ##2 ;
19797   {
19798     \int_compare:nNnTF
19799     { \_fp_array_count:n {##1} } = { \_fp_array_count:n {##2} }
19800     { \_fp_tuple_mapthread_o:nww { \_fp_binary_type_o:Nww #1 } }
19801     { \_fp_invalid_operation_o:nww #1 }
19802     \s_fp_tuple \_fp_tuple_chk:w {##1} ;
19803     \s_fp_tuple \_fp_tuple_chk:w {##2} ;
19804   }
19805 }
19806 \_fp_tmp:w +
19807 \_fp_tmp:w -

```

(End definition for `_fp_tuple+_tuple_o:ww` and `_fp_tuple-_tuple_o:ww`.)

```

19808 </package>

```

32 l3fp-extended implementation

```

19809 (*package)
19810 <@@=fp>

```

32.1 Description of fixed point numbers

This module provides a few functions to manipulate positive floating point numbers with extended precision (24 digits), but mostly provides functions for fixed-point numbers with this precision (24 digits). Those are used in the computation of Taylor series for the logarithm, exponential, and trigonometric functions. Since we eventually only care about the 16 first digits of the final result, some of the calculations are not performed with the full 24-digit precision. In other words, the last two blocks of each fixed point number may be wrong as long as the error is small enough to be rounded away when converting back to a floating point number. The fixed point numbers are expressed as

$$\{\langle a_1 \rangle\} \{\langle a_2 \rangle\} \{\langle a_3 \rangle\} \{\langle a_4 \rangle\} \{\langle a_5 \rangle\} \{\langle a_6 \rangle\} ;$$

where each $\langle a_i \rangle$ is exactly 4 digits (ranging from 0000 to 9999), except $\langle a_1 \rangle$, which may be any “not-too-large” non-negative integer, with or without leading zeros. Here, “not-too-large” depends on the specific function (see the corresponding comments for details). Checking for overflow is the responsibility of the code calling those functions. The fixed point number a corresponding to the representation above is $a = \sum_{i=1}^6 \langle a_i \rangle \cdot 10^{-4i}$.

Most functions we define here have the form

```
\__fp_fixed_⟨calculation⟩:wnn ⟨operand1⟩ ; ⟨operand2⟩ ; {⟨continuation⟩}
```

They perform the $\langle \text{calculation} \rangle$ on the two $\langle \text{operands} \rangle$, then feed the result (6 brace groups followed by a semicolon) to the $\langle \text{continuation} \rangle$, responsible for the next step of the calculation. Some functions only accept an N-type $\langle \text{continuation} \rangle$. This allows constructions such as

```
\__fp_fixed_add:wnn ⟨X1⟩ ; ⟨X2⟩ ;
\__fp_fixed_mul:wnn ⟨X3⟩ ;
\__fp_fixed_add:wnn ⟨X4⟩ ;
```

to compute $(X_1 + X_2) \cdot X_3 + X_4$. This turns out to be very appropriate for computing continued fractions and Taylor series.

At the end of the calculation, the result is turned back to a floating point number using `__fp_fixed_to_float_o:wN`. This function has to change the exponent of the floating point number: it must be used after starting an integer expression for the overall exponent of the result.

32.2 Helpers for numbers with extended precision

`\c__fp_one_fixed_t1` The fixed-point number 1, used in `l3fp-expo`.

```
19811 \tl_const:Nn \c__fp_one_fixed_t1
19812 { {10000} {0000} {0000} {0000} {0000} {0000} ; }
```

(End definition for `\c__fp_one_fixed_t1`.)

`__fp_fixed_continue:wn` This function simply calls the next function.

```
19813 \cs_new:Npn \__fp_fixed_continue:wn #1; #2 { #2 #1; }
```

(End definition for `__fp_fixed_continue:wn`.)

`__fp_fixed_add_one:wn`

`__fp_fixed_add_one:wn <a> ; <continuation>`

This function adds 1 to the fixed point $\langle a \rangle$, by changing a_1 to $10000 + a_1$, then calls the $\langle continuation \rangle$. This requires $a_1 + 10000 < 2^{31}$.

```
19814 \cs_new:Npn \__fp_fixed_add_one:wn #1#2; #3
19815 {
19816   \exp_after:wn #3 \exp_after:wn
19817   { \int_value:w \__fp_int_eval:w \c__fp_myriad_int + #1 } #2 ;
19818 }
```

(End definition for `__fp_fixed_add_one:wn`.)

`__fp_fixed_div_myriad:wn`

Divide a fixed point number by 10000. This is a little bit more subtle than just removing the last group and adding a leading group of zeros: the first group #1 may have any number of digits, and we must split #1 into the new first group and a second group of exactly 4 digits. The choice of shifts allows #1 to be in the range $[0, 5 \cdot 10^8 - 1]$.

```
19819 \cs_new:Npn \__fp_fixed_div_myriad:wn #1#2#3#4#5#6;
19820 {
19821   \exp_after:wn \__fp_fixed_mul_after:wnn
19822   \int_value:w \__fp_int_eval:w \c__fp_leading_shift_int
19823   \exp_after:wn \__fp_pack:NNNNNw
19824   \int_value:w \__fp_int_eval:w \c__fp_trailing_shift_int
19825   + #1 ; {#2}{#3}{#4}{#5};
19826 }
```

(End definition for `__fp_fixed_div_myriad:wn`.)

`__fp_fixed_mul_after:wnn`

The fixed point operations which involve multiplication end by calling this auxiliary. It braces the last block of digits, and places the $\langle continuation \rangle$ #3 in front.

```
19827 \cs_new:Npn \__fp_fixed_mul_after:wnn #1; #2; #3 { #3 {#1} #2; }
```

(End definition for `__fp_fixed_mul_after:wnn`.)

32.3 Multiplying a fixed point number by a short one

`__fp_fixed_mul_short:wnn`

```
\__fp_fixed_mul_short:wnn
{<a1>} {<a2>} {<a3>} {<a4>} {<a5>} {<a6>} ;
{<b0>} {<b1>} {<b2>} ; {<continuation>}
```

Computes the product $c = ab$ of $a = \sum_i \langle a_i \rangle 10^{-4i}$ and $b = \sum_i \langle b_i \rangle 10^{-4i}$, rounds it to the closest multiple of 10^{-24} , and leaves $\langle continuation \rangle$ $\{ \langle c_1 \rangle \} \dots \{ \langle c_6 \rangle \}$; in the input stream, where each of the $\langle c_i \rangle$ are blocks of 4 digits, except $\langle c_1 \rangle$, which is any \TeX integer. Note that indices for $\langle b \rangle$ start at 0: for instance a second operand of $\{0001\}\{0000\}\{0000\}$ leaves the first operand unchanged (rather than dividing it by 10^4 , as `__fp_fixed_mul:wnn` would).

```
19828 \cs_new:Npn \__fp_fixed_mul_short:wnn #1#2#3#4#5#6; #7#8#9;
19829 {
19830   \exp_after:wn \__fp_fixed_mul_after:wnn
19831   \int_value:w \__fp_int_eval:w \c__fp_leading_shift_int
19832   + #1*#7
19833   \exp_after:wn \__fp_pack:NNNNNw
19834   \int_value:w \__fp_int_eval:w \c__fp_middle_shift_int
19835   + #1*#8 + #2*#7
19836   \exp_after:wn \__fp_pack:NNNNNw
19837   \int_value:w \__fp_int_eval:w \c__fp_middle_shift_int
```



```

19838      + #1*#9 + #2*#8 + #3*#7
19839      \exp_after:wN \__fp_pack:NNNNNw
19840      \int_value:w \__fp_int_eval:w \c__fp_middle_shift_int
19841      + #2*#9 + #3*#8 + #4*#7
19842      \exp_after:wN \__fp_pack:NNNNNw
19843      \int_value:w \__fp_int_eval:w \c__fp_middle_shift_int
19844      + #3*#9 + #4*#8 + #5*#7
19845      \exp_after:wN \__fp_pack:NNNNNw
19846      \int_value:w \__fp_int_eval:w \c__fp_trailing_shift_int
19847      + #4*#9 + #5*#8 + #6*#7
19848      + ( #5*#9 + #6*#8 + #6*#9 / \c__fp_myriad_int )
19849      / \c__fp_myriad_int ; ;
19850  }

```

(End definition for __fp_fixed_mul_short:wN.)

32.4 Dividing a fixed point number by a small integer

```

\__fp_fixed_div_int:wwN \__fp_fixed_div_int:wwN <a> ; <n> ; <continuation>

```

Divides the fixed point number $\langle a \rangle$ by the (small) integer $0 < \langle n \rangle < 10^4$ and feeds the result to the $\langle continuation \rangle$. There is no bound on a_1 .

```

\__fp_fixed_div_int_auxi:wN

```

The arguments of the **i** auxiliary are 1: one of the a_i , 2: n , 3: the **ii** or the **iii** auxiliary. It computes a (somewhat tight) lower bound Q_i for the ratio a_i/n .

```

\__fp_fixed_div_int_auxii:wN

```

The **ii** auxiliary receives Q_i , n , and a_i as arguments. It adds Q_i to a surrounding integer expression, and starts a new one with the initial value 9999, which ensures that the result of this expression has 5 digits. The auxiliary also computes $a_i - n \cdot Q_i$, placing the result in front of the 4 digits of a_{i+1} . The resulting $a'_{i+1} = 10^4(a_i - n \cdot Q_i) + a_{i+1}$ serves as the first argument for a new call to the **i** auxiliary.

```

\__fp_fixed_div_int_pack:Nw

```

When the **iii** auxiliary is called, the situation looks like this:

```

\__fp_fixed_div_int_after:Nw <continuation>
-1 + Q1
\__fp_fixed_div_int_pack:Nw 9999 + Q2
\__fp_fixed_div_int_pack:Nw 9999 + Q3
\__fp_fixed_div_int_pack:Nw 9999 + Q4
\__fp_fixed_div_int_pack:Nw 9999 + Q5
\__fp_fixed_div_int_pack:Nw 9999
\__fp_fixed_div_int_auxii:wN Q6 ; {<n>} {<a6

```

where expansion is happening from the last line up. The **iii** auxiliary adds $Q_6 + 2 \simeq a_6/n + 1$ to the last 9999, giving the integer closest to $10000 + a_6/n$.

Each **pack** auxiliary receives 5 digits followed by a semicolon. The first digit is added as a carry to the integer expression above, and the 4 other digits are braced. Each call to the **pack** auxiliary thus produces one brace group. The last brace group is produced by the **after** auxiliary, which places the $\langle continuation \rangle$ as appropriate.

```

19851 \cs_new:Npn \__fp_fixed_div_int:wwN #1#2#3#4#5#6 ; #7 ; #8
19852 {
19853   \exp_after:wN \__fp_fixed_div_int_after:Nw
19854   \exp_after:wN #8
19855   \int_value:w \__fp_int_eval:w - 1
19856   \__fp_fixed_div_int:wwN
19857   #1; {#7} \__fp_fixed_div_int_auxi:wN

```

```

19858      #2; {#7} \__fp_fixed_div_int_auxi:wnn
19859      #3; {#7} \__fp_fixed_div_int_auxi:wnn
19860      #4; {#7} \__fp_fixed_div_int_auxi:wnn
19861      #5; {#7} \__fp_fixed_div_int_auxi:wnn
19862      #6; {#7} \__fp_fixed_div_int_auxii:wnn ;
19863    }
19864 \cs_new:Npn \__fp_fixed_div_int:wnN #1; #2 #3
19865 {
19866   \exp_after:wN #3
19867   \int_value:w \__fp_int_eval:w #1 / #2 - 1 ;
19868   {#2}
19869   {#1}
19870 }
19871 \cs_new:Npn \__fp_fixed_div_int_auxi:wnn #1; #2 #3
19872 {
19873   + #1
19874   \exp_after:wN \__fp_fixed_div_int_pack:Nw
19875   \int_value:w \__fp_int_eval:w 9999
19876   \exp_after:wN \__fp_fixed_div_int:wnN
19877   \int_value:w \__fp_int_eval:w #3 - #1*#2 \__fp_int_eval_end:
19878 }
19879 \cs_new:Npn \__fp_fixed_div_int_auxii:wnn #1; #2 #3 { + #1 + 2 ; }
19880 \cs_new:Npn \__fp_fixed_div_int_pack:Nw #1 #2; { + #1; {#2} }
19881 \cs_new:Npn \__fp_fixed_div_int_after:Nw #1 #2; { #1 {#2} }

```

(End definition for __fp_fixed_div_int:wnN and others.)

32.5 Adding and subtracting fixed points

`__fp_fixed_add:wnn` `__fp_fixed_add:wnn <a> ; ; {<continuation>}`
`__fp_fixed_sub:wnn` Computes $a + b$ (resp. $a - b$) and feeds the result to the $\langle continuation \rangle$. This function
`__fp_fixed_add:Nnnnnwnn` requires $0 \leq a_1, b_1 \leq 114748$, its result must be positive (this happens automatically for
`__fp_fixed_add:nnNnnwnn` addition) and its first group must have at most 5 digits: $(a \pm b)_1 < 100000$. The two
`__fp_fixed_add_pack:NNNNNwn` functions only differ by a sign, hence use a common auxiliary. It would be nice to grab
`__fp_fixed_add_after:NNNNNwn` the 12 brace groups in one go; only 9 parameters are allowed. Start by grabbing the sign,
 a_1, \dots, a_4 , the rest of a , and b_1 and b_2 . The second auxiliary receives the rest of a , the
sign multiplying b , the rest of b , and the $\langle continuation \rangle$ as arguments. After going down
through the various level, we go back up, packing digits and bringing the $\langle continuation \rangle$
(#8, then #7) from the end of the argument list to its start.

```

19882 \cs_new:Npn \__fp_fixed_add:wnn { \__fp_fixed_add:Nnnnnwnn + }
19883 \cs_new:Npn \__fp_fixed_sub:wnn { \__fp_fixed_add:Nnnnnwnn - }
19884 \cs_new:Npn \__fp_fixed_add:Nnnnnwnn #1 #2#3#4#5 #6; #7#8
19885 {
19886   \exp_after:wN \__fp_fixed_add_after:NNNNNwn
19887   \int_value:w \__fp_int_eval:w 9 9999 9998 + #2#3 #1 #7#8
19888   \exp_after:wN \__fp_fixed_add_pack:NNNNNwn
19889   \int_value:w \__fp_int_eval:w 1 9999 9998 + #4#5
19890   \__fp_fixed_add:nnNnnwn #6 #1
19891 }
19892 \cs_new:Npn \__fp_fixed_add:nnNnnwn #1#2 #3 #4#5 #6#7 ; #8
19893 {
19894   #3 #4#5
19895   \exp_after:wN \__fp_fixed_add_pack:NNNNNwn

```

```

19896 \int_value:w \__fp_int_eval:w 2 0000 0000 #3 #6#7 + #1#2 ; {#8} ;
19897 }
19898 \cs_new:Npn \__fp_fixed_add_pack:NNNNNwn #1 #2#3#4#5 #6; #7
19899 { + #1 ; {#7} {#2#3#4#5} {#6} }
19900 \cs_new:Npn \__fp_fixed_add_after:NNNNNwn 1 #1 #2#3#4#5 #6; #7
19901 { #7 {#1#2#3#4#5} {#6} }

```

(End definition for `__fp_fixed_add:wnn` and others.)

32.6 Multiplying fixed points

```

\__fp_fixed_mul:wnn
\__fp_fixed_mul:nnnnnnnw

```

`__fp_fixed_mul:wnn` $\langle a \rangle$; $\langle b \rangle$; $\{\langle continuation \rangle\}$

Computes $a \times b$ and feeds the result to $\langle continuation \rangle$. This function requires $0 \leq a_1, b_1 < 10000$. Once more, we need to play around the limit of 9 arguments for $\text{T}_{\text{E}}\text{X}$ macros. Note that we don't need to obtain an exact rounding, contrarily to the `*` operator, so things could be harder. We wish to perform carries in

$$\begin{aligned}
a \times b = & a_1 \cdot b_1 \cdot 10^{-8} \\
& + (a_1 \cdot b_2 + a_2 \cdot b_1) \cdot 10^{-12} \\
& + (a_1 \cdot b_3 + a_2 \cdot b_2 + a_3 \cdot b_1) \cdot 10^{-16} \\
& + (a_1 \cdot b_4 + a_2 \cdot b_3 + a_3 \cdot b_2 + a_4 \cdot b_1) \cdot 10^{-20} \\
& + \left(a_2 \cdot b_4 + a_3 \cdot b_3 + a_4 \cdot b_2 \right. \\
& \quad \left. + \frac{a_3 \cdot b_4 + a_4 \cdot b_3 + a_1 \cdot b_6 + a_2 \cdot b_5 + a_5 \cdot b_2 + a_6 \cdot b_1}{10^4} \right. \\
& \quad \left. + a_1 \cdot b_5 + a_5 \cdot b_1 \right) \cdot 10^{-24} + O(10^{-24}),
\end{aligned}$$

where the $O(10^{-24})$ stands for terms which are at most $5 \cdot 10^{-24}$; ignoring those leads to an error of at most 5 ulp. Note how the first 15 terms only depend on a_1, \dots, a_4 and b_1, \dots, b_4 , while the last 6 terms only depend on a_1, a_2, a_5, a_6 , and the corresponding parts of b . Hence, the first function grabs a_1, \dots, a_4 , the rest of a , and b_1, \dots, b_4 , and writes the 15 first terms of the expression, including a left parenthesis for the fraction. The `i` auxiliary receives $a_5, a_6, b_1, b_2, a_1, a_2, b_5, b_6$ and finally the $\langle continuation \rangle$ as arguments. It writes the end of the expression, including the right parenthesis and the denominator of the fraction. The $\langle continuation \rangle$ is finally placed in front of the 6 brace groups by `__fp_fixed_mul_after:wnn`.

```

19902 \cs_new:Npn \__fp_fixed_mul:wnn #1#2#3#4 #5; #6#7#8#9
19903 {
19904   \exp_after:wN \__fp_fixed_mul_after:wnn
19905   \int_value:w \__fp_int_eval:w \c__fp_leading_shift_int
19906   \exp_after:wN \__fp_pack:NNNNNw
19907   \int_value:w \__fp_int_eval:w \c__fp_middle_shift_int
19908   + #1*#6
19909   \exp_after:wN \__fp_pack:NNNNNw
19910   \int_value:w \__fp_int_eval:w \c__fp_middle_shift_int
19911   + #1*#7 + #2*#6
19912   \exp_after:wN \__fp_pack:NNNNNw
19913   \int_value:w \__fp_int_eval:w \c__fp_middle_shift_int
19914   + #1*#8 + #2*#7 + #3*#6
19915   \exp_after:wN \__fp_pack:NNNNNw

```

```

19916      \int_value:w \__fp_int_eval:w \c__fp_middle_shift_int
19917      + #1*#9 + #2*#8 + #3*#7 + #4*#6
19918      \exp_after:wN \__fp_pack:NNNNNw
19919      \int_value:w \__fp_int_eval:w \c__fp_trailing_shift_int
19920      + #2*#9 + #3*#8 + #4*#7
19921      + ( #3*#9 + #4*#8
19922      + \__fp_fixed_mul:nnnnnnnw #5 {#6}{#7} {#1}{#2}
19923      }
19924 \cs_new:Npn \__fp_fixed_mul:nnnnnnnw #1#2 #3#4 #5#6 #7#8 ;
19925 {
19926     #1*#4 + #2*#3 + #5*#8 + #6*#7 ) / \c__fp_myriad_int
19927     + #1*#3 + #5*#7 ; ;
19928 }

```

(End definition for `__fp_fixed_mul:wwn` and `__fp_fixed_mul:nnnnnnnw`.)

32.7 Combining product and sum of fixed points

```

\__fp_fixed_mul_add:wwwn
\__fp_fixed_mul_sub_back:wwwn
\__fp_fixed_one_minus_mul:wwn

```

Sometimes called FMA (fused multiply-add), these functions compute $a \times b + c$, $c - a \times b$, and $1 - a \times b$ and feed the result to the `\langle continuation \rangle`. Those functions require $0 \leq a_1, b_1, c_1 \leq 10000$. Since those functions are at the heart of the computation of Taylor expansions, we over-optimize them a bit, and in particular we do not factor out the common parts of the three functions.

For definiteness, consider the task of computing $a \times b + c$. We perform carries in

$$\begin{aligned}
 a \times b + c = & (a_1 \cdot b_1 + c_1 c_2) \cdot 10^{-8} \\
 & + (a_1 \cdot b_2 + a_2 \cdot b_1) \cdot 10^{-12} \\
 & + (a_1 \cdot b_3 + a_2 \cdot b_2 + a_3 \cdot b_1 + c_3 c_4) \cdot 10^{-16} \\
 & + (a_1 \cdot b_4 + a_2 \cdot b_3 + a_3 \cdot b_2 + a_4 \cdot b_1) \cdot 10^{-20} \\
 & + \left(a_2 \cdot b_4 + a_3 \cdot b_3 + a_4 \cdot b_2 \right. \\
 & \quad \left. + \frac{a_3 \cdot b_4 + a_4 \cdot b_3 + a_1 \cdot b_6 + a_2 \cdot b_5 + a_5 \cdot b_2 + a_6 \cdot b_1}{10^4} \right. \\
 & \quad \left. + a_1 \cdot b_5 + a_5 \cdot b_1 + c_5 c_6 \right) \cdot 10^{-24} + O(10^{-24}),
 \end{aligned}$$

where $c_1 c_2$, $c_3 c_4$, $c_5 c_6$ denote the 8-digit number obtained by juxtaposing the two blocks of digits of c , and \cdot denotes multiplication. The task is obviously tough because we have 18 brace groups in front of us.

Each of the three function starts the first two levels (the first, corresponding to 10^{-4} , is empty), with $c_1 c_2$ in the first level, calls the `i` auxiliary with arguments described later, and adds a trailing $+ c_5 c_6 ; \{\langle continuation \rangle\}$; . The $+ c_5 c_6$ piece, which is omitted for `__fp_fixed_one_minus_mul:wwn`, is taken in the integer expression for the 10^{-24} level.

```

19929 \cs_new:Npn \__fp_fixed_mul_add:wwwn #1; #2; #3#4#5#6#7#8;
19930 {
19931     \exp_after:wN \__fp_fixed_mul_after:wwn
19932     \int_value:w \__fp_int_eval:w \c__fp_big_leading_shift_int
19933     \exp_after:wN \__fp_pack_big:NNNNNNw

```

```

19934      \int_value:w \_fp_int_eval:w \c\_fp_big_middle_shift_int + #3 #4
19935      \_fp_fixed_mul_add:Nwnnnwnnn +
19936      + #5 #6 ; #2 ; #1 ; #2 ; +
19937      + #7 #8 ; ;
19938    }
19939 \cs_new:Npn \_fp_fixed_mul_sub_back:wwn #1; #2; #3#4#5#6#7#8;
19940 {
19941   \exp_after:wN \_fp_fixed_mul_after:wwn
19942   \int_value:w \_fp_int_eval:w \c\_fp_big_leading_shift_int
19943   \exp_after:wN \_fp_pack_big:NNNNNNw
19944   \int_value:w \_fp_int_eval:w \c\_fp_big_middle_shift_int + #3 #4
19945   \_fp_fixed_mul_add:Nwnnnwnnn -
19946   + #5 #6 ; #2 ; #1 ; #2 ; -
19947   + #7 #8 ; ;
19948 }
19949 \cs_new:Npn \_fp_fixed_one_minus_mul:wwn #1; #2;
19950 {
19951   \exp_after:wN \_fp_fixed_mul_after:wwn
19952   \int_value:w \_fp_int_eval:w \c\_fp_big_leading_shift_int
19953   \exp_after:wN \_fp_pack_big:NNNNNNw
19954   \int_value:w \_fp_int_eval:w \c\_fp_big_middle_shift_int +
19955   1 0000 0000
19956   \_fp_fixed_mul_add:Nwnnnwnnn -
19957   ; #2 ; #1 ; #2 ; -
19958   ; ;
19959 }

```

(End definition for `_fp_fixed_mul_add:wwn`, `_fp_fixed_mul_sub_back:wwn`, and `_fp_fixed_mul_one_minus_mul:wwn`.)

```

\_fp_fixed_mul_add:Nwnnnwnnn
      \_fp_fixed_mul_add:Nwnnnwnnn <op> + <c3> <c4> ;
      <b> ; <a> ; <b> ; <op>
      + <c5> <c6> ;

```

Here, `<op>` is either `+` or `-`. Arguments `#3`, `#4`, `#5` are $\langle b_1 \rangle$, $\langle b_2 \rangle$, $\langle b_3 \rangle$; arguments `#7`, `#8`, `#9` are $\langle a_1 \rangle$, $\langle a_2 \rangle$, $\langle a_3 \rangle$. We can build three levels: $a_1 \cdot b_1$ for 10^{-8} , $(a_1 \cdot b_2 + a_2 \cdot b_1)$ for 10^{-12} , and $(a_1 \cdot b_3 + a_2 \cdot b_2 + a_3 \cdot b_1 + c_3 c_4)$ for 10^{-16} . The a - b products use the sign `#1`. Note that `#2` is empty for `_fp_fixed_one_minus_mul:wwn`. We call the `ii` auxiliary for levels 10^{-20} and 10^{-24} , keeping the pieces of $\langle a \rangle$ we've read, but not $\langle b \rangle$, since there is another copy later in the input stream.

```

19960 \cs_new:Npn \_fp_fixed_mul_add:Nwnnnwnnn #1 #2; #3#4#5#6; #7#8#9
19961 {
19962   #1 #7*#3
19963   \exp_after:wN \_fp_pack_big:NNNNNNw
19964   \int_value:w \_fp_int_eval:w \c\_fp_big_middle_shift_int
19965   #1 #7*#4 #1 #8*#3
19966   \exp_after:wN \_fp_pack_big:NNNNNNw
19967   \int_value:w \_fp_int_eval:w \c\_fp_big_middle_shift_int
19968   #1 #7*#5 #1 #8*#4 #1 #9*#3 #2
19969   \exp_after:wN \_fp_pack_big:NNNNNNw
19970   \int_value:w \_fp_int_eval:w \c\_fp_big_middle_shift_int
19971   #1 \_fp_fixed_mul_add:nnnnwnnn {#7}{#8}{#9}
19972 }

```

(End definition for `_fp_fixed_mul_add:Nwnnnwnnn`.)

_fp_fixed_mul_add:nnnnwnnnn

_fp_fixed_mul_add:nnnnwnnnn $\langle a \rangle$; $\langle b \rangle$; $\langle op \rangle$
 $+ \langle c_5 \rangle \langle c_6 \rangle$;

Level 10^{-20} is $(a_1 \cdot b_4 + a_2 \cdot b_3 + a_3 \cdot b_2 + a_4 \cdot b_1)$, multiplied by the sign, which was inserted by the *i* auxiliary. Then we prepare level 10^{-24} . We don't have access to all parts of $\langle a \rangle$ and $\langle b \rangle$ needed to make all products. Instead, we prepare the partial expressions

$$b_1 + a_4 \cdot b_2 + a_3 \cdot b_3 + a_2 \cdot b_4 + a_1$$

$$b_2 + a_4 \cdot b_3 + a_3 \cdot b_4 + a_2.$$

Obviously, those expressions make no mathematical sense: we complete them with $a_5 \cdot$ and $\cdot b_5$, and with $a_6 \cdot b_1 + a_5 \cdot$ and $\cdot b_5 + a_1 \cdot b_6$, and of course with the trailing $+ c_5 c_6$. To do all this, we keep a_1, a_5, a_6 , and the corresponding pieces of $\langle b \rangle$.

```

19973 \cs_new:Npn \_fp\_fixed\_mul\_add:nnnnwnnnn #1#2#3#4#5; #6#7#8#9
19974 {
19975   ( #1*#9 + #2*#8 + #3*#7 + #4*#6 )
19976   \exp_after:wN \_fp\_pack\_big:NNNNNNw
19977   \int_value:w \_fp\_int\_eval:w \c\_fp\_big\_trailing\_shift\_int
19978   \_fp\_fixed\_mul\_add:nnnnwnnwN
19979   { #6 + #4*#7 + #3*#8 + #2*#9 + #1 }
19980   { #7 + #4*#8 + #3*#9 + #2 }
19981   {#1} #5;
19982   {#6}
19983 }

```

(End definition for _fp_fixed_mul_add:nnnnwnnnn.)

_fp_fixed_mul_add:nnnnwnnwN

_fp_fixed_mul_add:nnnnwnnwN $\{\langle partial_1 \rangle\} \{\langle partial_2 \rangle\}$
 $\{\langle a_1 \rangle\} \{\langle a_5 \rangle\} \{\langle a_6 \rangle\}$; $\{\langle b_1 \rangle\} \{\langle b_5 \rangle\} \{\langle b_6 \rangle\}$;
 $\langle op \rangle + \langle c_5 \rangle \langle c_6 \rangle$;

Complete the $\langle partial_1 \rangle$ and $\langle partial_2 \rangle$ expressions as explained for the *ii* auxiliary. The second one is divided by 10000: this is the carry from level 10^{-28} . The trailing $+ c_5 c_6$ is taken into the expression for level 10^{-24} . Note that the total of level 10^{-24} is in the interval $[-5 \cdot 10^8, 6 \cdot 10^8]$ (give or take a couple of 10000), hence adding it to the shift gives a 10-digit number, as expected by the packing auxiliaries. See *l3fp-aux* for the definition of the shifts and packing auxiliaries.

```

19984 \cs_new:Npn \_fp\_fixed\_mul\_add:nnnnwnnwN #1#2 #3#4#5; #6#7#8; #9
19985 {
19986   #9 (#4* #1 *#7)
19987   #9 (#5*#6+#4* #2 *#7+#3*#8) / \c\_fp\_myriad\_int
19988 }

```

(End definition for _fp_fixed_mul_add:nnnnwnnwN.)

32.8 Extended-precision floating point numbers

In this section we manipulate floating point numbers with roughly 24 significant figures (“extended-precision” numbers, in short, “ep”), which take the form of an integer exponent, followed by a comma, then six groups of digits, ending with a semicolon. The first group of digit may be any non-negative integer, while other groups of digits have 4 digits. In other words, an extended-precision number is an exponent ending in a comma, then a fixed point number. The corresponding value is $0.\langle digits \rangle \cdot 10^{\langle exponent \rangle}$. This convention differs from floating points.

```

    \__fp_ep_to_fixed:wwn
    \__fp_ep_to_fixed_auxi:www
    \__fp_ep_to_fixed_auxii:nnnnnnnnwn

```

Converts an extended-precision number with an exponent at most 4 and a first block less than 10^8 to a fixed point number whose first block has 12 digits, hopefully starting with many zeros.

```

19989 \cs_new:Npn \__fp_ep_to_fixed:wwn #1,#2
19990 {
19991     \exp_after:wN \__fp_ep_to_fixed_auxi:www
19992     \int_value:w \__fp_int_eval:w 1 0000 0000 + #2 \exp_after:wN ;
19993     \exp:w \exp_end_continue_f:w
19994     \prg_replicate:nn { 4 - \int_max:nn {#1} { -32 } } { 0 } ;
19995 }
19996 \cs_new:Npn \__fp_ep_to_fixed_auxi:www 1#1; #2; #3#4#5#6#7;
19997 {
19998     \__fp_pack_eight:wNNNNNNNN
19999     \__fp_pack_twice_four:wNNNNNNNN
20000     \__fp_pack_twice_four:wNNNNNNNN
20001     \__fp_pack_twice_four:wNNNNNNNN
20002     \__fp_ep_to_fixed_auxii:nnnnnnnnwn ;
20003     #2 #1#3#4#5#6#7 0000 !
20004 }
20005 \cs_new:Npn \__fp_ep_to_fixed_auxii:nnnnnnnnwn #1#2#3#4#5#6#7; #8! #9
20006 { #9 {#1#2}{#3}{#4}{#5}{#6}{#7}; }

```

(End definition for __fp_ep_to_fixed:wwn, __fp_ep_to_fixed_auxi:www, and __fp_ep_to_fixed_auxii:nnnnnnnnwn.)

```

    \__fp_ep_to_ep:wwN
    \__fp_ep_to_ep_loop:N
    \__fp_ep_to_ep_end:www
    \__fp_ep_to_ep_zero:ww

```

Normalize an extended-precision number. More precisely, leading zeros are removed from the mantissa of the argument, decreasing its exponent as appropriate. Then the digits are packed into 6 groups of 4 (discarding any remaining digit, not rounding). Finally, the continuation #8 is placed before the resulting exponent-mantissa pair. The input exponent may in fact be given as an integer expression. The loop auxiliary grabs a digit: if it is 0, decrement the exponent and continue looping, and otherwise call the end auxiliary, which places all digits in the right order (the digit that was not 0, and any remaining digits), followed by some 0, then packs them up neatly in $3 \times 2 = 6$ blocks of four. At the end of the day, remove with __fp_use_i:ww any digit that did not make it in the final mantissa (typically only zeros, unless the original first block has more than 4 digits).

```

20007 \cs_new:Npn \__fp_ep_to_ep:wwN #1,#2#3#4#5#6#7; #8
20008 {
20009     \exp_after:wN #8
20010     \int_value:w \__fp_int_eval:w #1 + 4
20011     \exp_after:wN \use_i:nn
20012     \exp_after:wN \__fp_ep_to_ep_loop:N
20013     \int_value:w \__fp_int_eval:w 1 0000 0000 + #2 \__fp_int_eval_end:
20014     #3#4#5#6#7 ; ; !
20015 }
20016 \cs_new:Npn \__fp_ep_to_ep_loop:N #1
20017 {
20018     \if_meaning:w 0 #1
20019     - 1
20020     \else:
20021         \__fp_ep_to_ep_end:www #1
20022     \fi:
20023     \__fp_ep_to_ep_loop:N

```

```

20024     }
20025 \cs_new:Npn \__fp_ep_to_ep_end:www
20026   #1 \fi: \__fp_ep_to_ep_loop:N #2; #3!
20027   {
20028     \fi:
20029     \if_meaning:w ; #1
20030       - 2 * \c_fp_max_exponent_int
20031     \__fp_ep_to_ep_zero:ww
20032     \fi:
20033     \__fp_pack_twice_four:wNNNNNNNN
20034     \__fp_pack_twice_four:wNNNNNNNN
20035     \__fp_pack_twice_four:wNNNNNNNN
20036     \__fp_use_i:ww , ;
20037     #1 #2 0000 0000 0000 0000 0000 0000 ;
20038   }
20039 \cs_new:Npn \__fp_ep_to_ep_zero:ww \fi: #1; #2; #3;
20040   { \fi: , {1000}{0000}{0000}{0000}{0000}{0000} ; }

```

(End definition for __fp_ep_to_ep:wwN and others.)

__fp_ep_compare:www
__fp_ep_compare_aux:www

In l3fp-trig we need to compare two extended-precision numbers. This is based on the same function for positive floating point numbers, with an extra test if comparing only 16 decimals is not enough to distinguish the numbers. Note that this function only works if the numbers are normalized so that their first block is in [1000,9999].

```

20041 \cs_new:Npn \__fp_ep_compare:www #1,#2#3#4#5#6#7;
20042   { \__fp_ep_compare_aux:www {#1}{#2}{#3}{#4}{#5}; #6#7; }
20043 \cs_new:Npn \__fp_ep_compare_aux:www #1;#2;#3,#4#5#6#7#8#9;
20044   {
20045     \if_case:w
20046       \__fp_compare_npos:nwnw #1; {#3}{#4}{#5}{#6}{#7}; \exp_stop_f:
20047       \if_int_compare:w #2 = #8#9 \exp_stop_f:
20048         0
20049       \else:
20050         \if_int_compare:w #2 < #8#9 - \fi: 1
20051       \fi:
20052     \or: 1
20053     \else: -1
20054     \fi:
20055   }

```

(End definition for __fp_ep_compare:www and __fp_ep_compare_aux:www.)

__fp_ep_mul:wwwN
__fp_ep_mul_raw:wwwN

Multiply two extended-precision numbers: first normalize them to avoid losing too much precision, then multiply the mantissas #2 and #4 as fixed point numbers, and sum the exponents #1 and #3. The result's first block is in [100,9999].

```

20056 \cs_new:Npn \__fp_ep_mul:wwwN #1,#2; #3,#4;
20057   {
20058     \__fp_ep_to_ep:wwN #3,#4;
20059     \__fp_fixed_continue:wn
20060     {
20061       \__fp_ep_to_ep:wwN #1,#2;
20062       \__fp_ep_mul_raw:wwwN
20063     }
20064     \__fp_fixed_continue:wn

```



```

20065     }
20066 \cs_new:Npn \__fp_ep_mul_raw:wwwN #1,#2; #3,#4; #5
20067     {
20068     \__fp_fixed_mul:wn #2; #4;
20069     { \exp_after:wN #5 \int_value:w \__fp_int_eval:w #1 + #3 , }
20070     }

```

(End definition for `__fp_ep_mul:wwwN` and `__fp_ep_mul_raw:wwwN`.)

32.9 Dividing extended-precision numbers

Divisions of extended-precision numbers are difficult to perform with exact rounding: the technique used in `l3fp-basics` for 16-digit floating point numbers does not generalize easily to 24-digit numbers. Thankfully, there is no need for exact rounding.

Let us call $\langle n \rangle$ the numerator and $\langle d \rangle$ the denominator. After a simple normalization step, we can assume that $\langle n \rangle \in [0.1, 1)$ and $\langle d \rangle \in [0.1, 1)$, and compute $\langle n \rangle / (10 \langle d \rangle) \in (0.01, 1)$. In terms of the 6 blocks of digits $\langle n_1 \rangle \cdots \langle n_6 \rangle$ and the 6 blocks $\langle d_1 \rangle \cdots \langle d_6 \rangle$, the condition translates to $\langle n_1 \rangle, \langle d_1 \rangle \in [1000, 9999]$.

We first find an integer estimate $a \simeq 10^8 / \langle d \rangle$ by computing

$$\alpha = \left\lceil \frac{10^9}{\langle d_1 \rangle + 1} \right\rceil$$

$$\beta = \left\lfloor \frac{10^9}{\langle d_1 \rangle} \right\rfloor$$

$$a = 10^3 \alpha + (\beta - \alpha) \cdot \left(10^3 - \left\lfloor \frac{\langle d_2 \rangle}{10} \right\rfloor \right) - 1250,$$

where $\left\lceil \cdot \right\rceil$ denotes ε -TEX's rounding division, which rounds ties away from zero. The idea is to interpolate between $10^3 \alpha$ and $10^3 \beta$ with a parameter $\langle d_2 \rangle / 10^4$, so that when $\langle d_2 \rangle = 0$ one gets $a = 10^3 \beta - 1250 \simeq 10^{12} / \langle d_1 \rangle \simeq 10^8 / \langle d \rangle$, while when $\langle d_2 \rangle = 9999$ one gets $a = 10^3 \alpha - 1250 \simeq 10^{12} / (\langle d_1 \rangle + 1) \simeq 10^8 / \langle d \rangle$. The shift by 1250 helps to ensure that a is an underestimate of the correct value. We shall prove that

$$1 - 1.755 \cdot 10^{-5} < \frac{\langle d \rangle a}{10^8} < 1.$$

We can then compute the inverse of $\langle d \rangle a / 10^8 = 1 - \epsilon$ using the relation $1 / (1 - \epsilon) \simeq (1 + \epsilon)(1 + \epsilon^2) + \epsilon^4$, which is correct up to a relative error of $\epsilon^5 < 1.6 \cdot 10^{-24}$. This allows us to find the desired ratio as

$$\frac{\langle n \rangle}{\langle d \rangle} = \frac{\langle n \rangle a}{10^8} ((1 + \epsilon)(1 + \epsilon^2) + \epsilon^4).$$

Let us prove the upper bound first (multiplied by 10^{15}). Note that $10^7 \langle d \rangle < 10^3 \langle d_1 \rangle + 10^{-1}(\langle d_2 \rangle + 1)$, and that ε -TEX's division $\left\lfloor \frac{\langle d_2 \rangle}{10} \right\rfloor$ underestimates $10^{-1}(\langle d_2 \rangle + 1)$ by 0.5 at

most, as can be checked for each possible last digit of $\langle d_2 \rangle$. Then,

$$10^7 \langle d \rangle a < \left(10^3 \langle d_1 \rangle + \left\lfloor \frac{\langle d_2 \rangle}{10} \right\rfloor + \frac{1}{2} \right) \left(\left(10^3 - \left\lfloor \frac{\langle d_2 \rangle}{10} \right\rfloor \right) \beta + \left\lfloor \frac{\langle d_2 \rangle}{10} \right\rfloor \alpha - 1250 \right) \quad (1)$$

$$< \left(10^3 \langle d_1 \rangle + \left\lfloor \frac{\langle d_2 \rangle}{10} \right\rfloor + \frac{1}{2} \right) \quad (2)$$

$$\left(\left(10^3 - \left\lfloor \frac{\langle d_2 \rangle}{10} \right\rfloor \right) \left(\frac{10^9}{\langle d_1 \rangle} + \frac{1}{2} \right) + \left\lfloor \frac{\langle d_2 \rangle}{10} \right\rfloor \left(\frac{10^9}{\langle d_1 \rangle + 1} + \frac{1}{2} \right) - 1250 \right) \quad (3)$$

$$< \left(10^3 \langle d_1 \rangle + \left\lfloor \frac{\langle d_2 \rangle}{10} \right\rfloor + \frac{1}{2} \right) \left(\frac{10^{12}}{\langle d_1 \rangle} - \left\lfloor \frac{\langle d_2 \rangle}{10} \right\rfloor \frac{10^9}{\langle d_1 \rangle (\langle d_1 \rangle + 1)} - 750 \right) \quad (4)$$

We recognize a quadratic polynomial in $[\langle d_2 \rangle / 10]$ with a negative leading coefficient: this polynomial is bounded above, according to $([\langle d_2 \rangle / 10] + a)(b - c[\langle d_2 \rangle / 10]) \leq (b + ca)^2 / (4c)$. Hence,

$$10^7 \langle d \rangle a < \frac{10^{15}}{\langle d_1 \rangle (\langle d_1 \rangle + 1)} \left(\langle d_1 \rangle + \frac{1}{2} + \frac{1}{4} 10^{-3} - \frac{3}{8} \cdot 10^{-9} \langle d_1 \rangle (\langle d_1 \rangle + 1) \right)^2$$

Since $\langle d_1 \rangle$ takes integer values within $[1000, 9999]$, it is a simple programming exercise to check that the squared expression is always less than $\langle d_1 \rangle (\langle d_1 \rangle + 1)$, hence $10^7 \langle d \rangle a < 10^{15}$. The upper bound is proven. We also find that $\frac{3}{8}$ can be replaced by slightly smaller numbers, but nothing less than $0.374563\dots$, and going back through the derivation of the upper bound, we find that 1250 is as small a shift as we can obtain without breaking the bound.

Now, the lower bound. The same computation as for the upper bound implies

$$10^7 \langle d \rangle a > \left(10^3 \langle d_1 \rangle + \left\lfloor \frac{\langle d_2 \rangle}{10} \right\rfloor - \frac{1}{2} \right) \left(\frac{10^{12}}{\langle d_1 \rangle} - \left\lfloor \frac{\langle d_2 \rangle}{10} \right\rfloor \frac{10^9}{\langle d_1 \rangle (\langle d_1 \rangle + 1)} - 1750 \right)$$

This time, we want to find the minimum of this quadratic polynomial. Since the leading coefficient is still negative, the minimum is reached for one of the extreme values $[y/10] = 0$ or $[y/10] = 100$, and we easily check the bound for those values.

We have proven that the algorithm gives us a precise enough answer. Incidentally, the upper bound that we derived tells us that $a < 10^8 / \langle d \rangle \leq 10^9$, hence we can compute a safely as a $\text{T}_{\text{E}}\text{X}$ integer, and even add 10^9 to it to ease grabbing of all the digits. The lower bound implies $10^8 - 1755 < a$, which we do not care about.

`__fp_ep_div:wwwn` Compute the ratio of two extended-precision numbers. The result is an extended-precision number whose first block lies in the range $[100, 9999]$, and is placed after the $\langle continuation \rangle$ once we are done. First normalize the inputs so that both first block lie in $[1000, 9999]$, then call `__fp_ep_div_esti:wwwn $\langle denominator \rangle$ $\langle numerator \rangle$` , responsible for estimating the inverse of the denominator.

```

20071 \cs_new:Npn \__fp_ep_div:wwwn #1,#2; #3,#4;
20072 {
20073   \__fp_ep_to_ep:wwN #1,#2;
20074   \__fp_fixed_continue:wn
20075   {
20076     \__fp_ep_to_ep:wwN #3,#4;
20077     \__fp_ep_div_esti:wwwn
20078   }
20079 }
```

(End definition for _fp_ep_div:wwwn.)

_fp_ep_div_esti:wwwn
_fp_ep_div_estii:wnnnwn
_fp_ep_div_estiii:NNNNwwn

The **esti** function evaluates $\alpha = 10^9 / (\langle d_1 \rangle + 1)$, which is used twice in the expression for a , and combines the exponents **#1** and **#4** (with a shift by 1 because we later compute $\langle n \rangle / (10 \langle d \rangle)$). Then the **estii** function evaluates $10^9 + a$, and puts the exponent **#2** after the continuation **#7**: from there on we can forget exponents and focus on the mantissa. The **estiii** function multiplies the denominator **#7** by $10^{-8}a$ (obtained as a split into the single digit **#1** and two blocks of 4 digits, **#2#3#4#5** and **#6**). The result $10^{-8}a \langle d \rangle = (1 - \epsilon)$, and a partially packed $10^{-9}a$ (as a block of four digits, and five individual digits, not packed by lack of available macro parameters here) are passed to **_fp_ep_div_epsilon:wnNNNNn**, which computes $10^{-9}a / (1 - \epsilon)$, that is, $1 / (10 \langle d \rangle)$ and we finally multiply this by the numerator **#8**.

```
20080 \cs_new:Npn \_fp_ep_div_esti:wwwn #1,#2#3; #4,
20081 {
20082   \exp_after:wN \_fp_ep_div_estii:wnnnwn
20083   \int_value:w \_fp_int_eval:w 10 0000 0000 / ( #2 + 1 )
20084   \exp_after:wN ;
20085   \int_value:w \_fp_int_eval:w #4 - #1 + 1 ,
20086   {#2} #3;
20087 }
20088 \cs_new:Npn \_fp_ep_div_estii:wnnnwn #1; #2,#3#4#5; #6; #7
20089 {
20090   \exp_after:wN \_fp_ep_div_estiii:NNNNwwn
20091   \int_value:w \_fp_int_eval:w 10 0000 0000 - 1750
20092   + #1 000 + (10 0000 0000 / #3 - #1) * (1000 - #4 / 10) ;
20093   {#3}{#4}#5; #6; { #7 #2, }
20094 }
20095 \cs_new:Npn \_fp_ep_div_estiii:NNNNwwn 1#1#2#3#4#5#6; #7;
20096 {
20097   \_fp_fixed_mul_short:wnn #7; {#1}{#2#3#4#5}{#6};
20098   \_fp_ep_div_epsilon:wnNNNNn {#1#2#3#4}#5#6
20099   \_fp_fixed_mul:wnn
20100 }
```

(End definition for _fp_ep_div_esti:wwwn, _fp_ep_div_estii:wnnnwn, and _fp_ep_div_estiii:NNNNwwn.)

_fp_ep_div_epsilon:wnNNNNn
_fp_ep_div_eps_pack:NNNNw
_fp_ep_div_epsii:wnNNNNn

The bounds shown above imply that the **epsi** function's first operand is $(1 - \epsilon)$ with $\epsilon \in [0, 1.755 \cdot 10^{-5}]$. The **epsi** function computes ϵ as $1 - (1 - \epsilon)$. Since $\epsilon < 10^{-4}$, its first block vanishes and there is no need to explicitly use **#1** (which is 9999). Then **epsii** evaluates $10^{-9}a / (1 - \epsilon)$ as $(1 + \epsilon^2)(1 + \epsilon)(10^{-9}a\epsilon) + 10^{-9}a$. Importantly, we compute $10^{-9}a\epsilon$ before multiplying it with the rest, rather than multiplying by ϵ and then $10^{-9}a$, as this second option loses more precision. Also, the combination of **short_mul** and **div_myriad** is both faster and more precise than a simple **mul**.

```
20101 \cs_new:Npn \_fp_ep_div_epsilon:wnNNNNn #1#2#3#4#5#6;
20102 {
20103   \exp_after:wN \_fp_ep_div_epsii:wnNNNNn
20104   \int_value:w \_fp_int_eval:w 1 9998 - #2
20105   \exp_after:wN \_fp_ep_div_eps_pack:NNNNw
20106   \int_value:w \_fp_int_eval:w 1 9999 9998 - #3#4
20107   \exp_after:wN \_fp_ep_div_eps_pack:NNNNw
20108   \int_value:w \_fp_int_eval:w 2 0000 0000 - #5#6 ; ;
20109 }
```

```

20110 \cs_new:Npn \__fp_ep_div_eps_pack:NNNNNw #1#2#3#4#5#6;
20111 { + #1 ; {#2#3#4#5} {#6} }
20112 \cs_new:Npn \__fp_ep_div_epsii:wnNNNNNn 1#1; #2; #3#4#5#6#7#8
20113 {
20114   \__fp_fixed_mul:wnn {0000}{#1}#2; {0000}{#1}#2;
20115   \__fp_fixed_add_one:wN
20116   \__fp_fixed_mul:wnn {10000} {#1} #2 ;
20117   {
20118     \__fp_fixed_mul_short:wnn {0000}{#1}#2; {#3}{#4#5#6#7}{#8000};
20119     \__fp_fixed_div_myriad:wn
20120     \__fp_fixed_mul:wnn
20121   }
20122   \__fp_fixed_add:wnn {#3}{#4#5#6#7}{#8000}{0000}{0000}{0000};
20123 }

```

(End definition for `__fp_ep_div_epsilon:wnNNNNNn`, `__fp_ep_div_eps_pack:NNNNNw`, and `__fp_ep_div_epsii:wnNNNNNn`.)

32.10 Inverse square root of extended precision numbers

The idea here is similar to division. Normalize the input, multiplying by powers of 100 until we have $x \in [0.01, 1)$. Then find an integer approximation $r \in [101, 1003]$ of $10^2/\sqrt{x}$, as the fixed point of iterations of the Newton method: essentially $r \mapsto (r + 10^8/(x_1 r))/2$, starting from a guess that optimizes the number of steps before convergence. In fact, just as there is a slight shift when computing divisions to ensure that some inequalities hold, we replace 10^8 by a slightly larger number which ensures that $r^2 x \geq 10^4$. This also causes $r \in [101, 1003]$. Another correction to the above is that the input is actually normalized to $[0.1, 1)$, and we use either 10^8 or 10^9 in the Newton method, depending on the parity of the exponent. Skipping those technical hurdles, once we have the approximation r , we set $y = 10^{-4}r^2x$ (or rather, the correct power of 10 to get $y \simeq 1$) and compute $y^{-1/2}$ through another application of Newton's method. This time, the starting value is $z = 1$, each step maps $z \mapsto z(1.5 - 0.5yz^2)$, and we perform a fixed number of steps. Our final result combines r with $y^{-1/2}$ as $x^{-1/2} = 10^{-2}ry^{-1/2}$.

```

\__fp_ep_isqrt:wnn
\__fp_ep_isqrt_aux:wnn
\__fp_ep_isqrt_auxii:wnnnwnn

```

First normalize the input, then check the parity of the exponent #1. If it is even, the result's exponent will be $-#1/2$, otherwise it will be $(#1 - 1)/2$ (except in the case where the input was an exact power of 100). The `auxii` function receives as #1 the result's exponent just computed, as #2 the starting value for the iteration giving r (the values 168 and 535 lead to the least number of iterations before convergence, on average), as #3 and #4 one empty argument and one 0, depending on the parity of the original exponent, as #5 and #6 the normalized mantissa (#5 $\in [1000, 9999]$), and as #7 the continuation. It sets up the iteration giving r : the `esti` function thus receives the initial two guesses #2 and 0, an approximation #5 of 10^4x (its first block of digits), and the empty/zero arguments #3 and #4, followed by the mantissa and an altered continuation where we have stored the result's exponent.

```

20124 \cs_new:Npn \__fp_ep_isqrt:wnn #1,#2;
20125 {
20126   \__fp_ep_to_ep:wnN #1,#2;
20127   \__fp_ep_isqrt_auxi:wnn
20128 }
20129 \cs_new:Npn \__fp_ep_isqrt_auxi:wnn #1,
20130 {

```

```

20131 \exp_after:wN \_fp_ep_isqrt_auxii:wwnnwn
20132 \int_value:w \_fp_int_eval:w
20133 \int_if_odd:nTF {#1}
20134 { (1 - #1) / 2 , 535 , { 0 } { } }
20135 { 1 - #1 / 2 , 168 , { } { 0 } }
20136 }
20137 \cs_new:Npn \_fp_ep_isqrt_auxii:wwnnwn #1, #2, #3#4 #5#6; #7
20138 {
20139 \_fp_ep_isqrt_esti:wwnnwn #2, 0, #5, {#3} {#4}
20140 {#5} #6 ; { #7 #1 , }
20141 }

```

(End definition for _fp_ep_isqrt:wn, _fp_ep_isqrt_aux:wn, and _fp_ep_isqrt_auxii:wwnnwn.)

```

\_fp_ep_isqrt_esti:wwnnwn
\_fp_ep_isqrt_estii:wwnnwn
\_fp_ep_isqrt_estiii:NNNNNwwn

```

If the last two approximations gave the same result, we are done: call the `esti` function to clean up. Otherwise, evaluate $(\langle prev \rangle + 1.005 \cdot 10^8 \text{ or } 9 / (\langle prev \rangle \cdot x)) / 2$, as the next approximation: omitting the 1.005 factor, this would be Newton's method. We can check by brute force that if `#4` is empty (the original exponent was even), the process computes an integer slightly larger than $100/\sqrt{x}$, while if `#4` is 0 (the original exponent was odd), the result is an integer slightly larger than $100/\sqrt{x/10}$. Once we are done, we evaluate $100r^2/2$ or $10r^2/2$ (when the exponent is even or odd, respectively) and feed that to `estiii`. This third auxiliary finds $y_{\text{even}}/2 = 10^{-4}r^2x/2$ or $y_{\text{odd}}/2 = 10^{-5}r^2x/2$ (again, depending on earlier parity). A simple program shows that $y \in [1, 1.0201]$. The number $y/2$ is fed to `_fp_ep_isqrt_epsilon:wn`, which computes $1/\sqrt{y}$, and we finally multiply the result by r .

```

20142 \cs_new:Npn \_fp_ep_isqrt_esti:wwnnwn #1, #2, #3, #4
20143 {
20144 \if_int_compare:w #1 = #2 \exp_stop_f:
20145 \exp_after:wN \_fp_ep_isqrt_estii:wwnnwn
20146 \fi:
20147 \exp_after:wN \_fp_ep_isqrt_esti:wwnnwn
20148 \int_value:w \_fp_int_eval:w
20149 (#1 + 1 0050 0000 #4 / (#1 * #3)) / 2 ,
20150 #1, #3, {#4}
20151 }
20152 \cs_new:Npn \_fp_ep_isqrt_estii:wwnnwn #1, #2, #3, #4#5
20153 {
20154 \exp_after:wN \_fp_ep_isqrt_estiii:NNNNNwwn
20155 \int_value:w \_fp_int_eval:w 1000 0000 + #2 * #2 #5 * 5
20156 \exp_after:wN , \int_value:w \_fp_int_eval:w 10000 + #2 ;
20157 }
20158 \cs_new:Npn \_fp_ep_isqrt_estiii:NNNNNwwn 1#1#2#3#4#5#6, 1#7#8; #9;
20159 {
20160 \_fp_fixed_mul_short:wn #9; {#1} {#2#3#4#5} {#600} ;
20161 \_fp_ep_isqrt_epsilon:wn
20162 \_fp_fixed_mul_short:wn {#7} {#80} {0000} ;
20163 }

```

(End definition for _fp_ep_isqrt_esti:wwnnwn, _fp_ep_isqrt_estii:wwnnwn, and _fp_ep_isqrt_estiii:NNNNNwwn.)

```

\_fp_ep_isqrt_epsilon:wn
\_fp_ep_isqrt_epsilonii:wn

```

Here, we receive a fixed point number $y/2$ with $y \in [1, 1.0201]$. Starting from $z = 1$ we iterate $z \mapsto z(3/2 - z^2y/2)$. In fact, we start from the first iteration $z = 3/2 - y/2$ to avoid useless multiplications. The `epsii` auxiliary receives z as `#1` and y as `#2`.

```

20164 \cs_new:Npn \__fp_ep_isqrt_epsi:wwN #1;
20165 {
20166   \__fp_fixed_sub:wwn {15000}{0000}{0000}{0000}{0000}{0000}; #1;
20167   \__fp_ep_isqrt_epsi:wwN #1;
20168   \__fp_ep_isqrt_epsi:wwN #1;
20169   \__fp_ep_isqrt_epsi:wwN #1;
20170 }
20171 \cs_new:Npn \__fp_ep_isqrt_epsi:wwN #1; #2;
20172 {
20173   \__fp_fixed_mul:wwn #1; #1;
20174   \__fp_fixed_mul_sub_back:wwwn #2;
20175   {15000}{0000}{0000}{0000}{0000}{0000};
20176   \__fp_fixed_mul:wwn #1;
20177 }

```

(End definition for __fp_ep_isqrt_epsi:wwN and __fp_ep_isqrt_epsi:wwN.)

32.11 Converting from fixed point to floating point

After computing Taylor series, we wish to convert the result from extended precision (with or without an exponent) to the public floating point format. The functions here should be called within an integer expression for the overall exponent of the floating point.

__fp_ep_to_float_o:wwN
 __fp_ep_inv_to_float_o:wwN

An extended-precision number is simply a comma-delimited exponent followed by a fixed point number. Leave the exponent in the current integer expression then convert the fixed point number.

```

20178 \cs_new:Npn \__fp_ep_to_float_o:wwN #1,
20179 { + \__fp_int_eval:w #1 \__fp_fixed_to_float_o:wwN }
20180 \cs_new:Npn \__fp_ep_inv_to_float_o:wwN #1,#2;
20181 {
20182   \__fp_ep_div:wwwn 1,{1000}{0000}{0000}{0000}{0000}{0000}; #1,#2;
20183   \__fp_ep_to_float_o:wwN
20184 }

```

(End definition for __fp_ep_to_float_o:wwN and __fp_ep_inv_to_float_o:wwN.)

__fp_fixed_inv_to_float_o:wwN

Another function which reduces to converting an extended precision number to a float.

```

20185 \cs_new:Npn \__fp_fixed_inv_to_float_o:wwN
20186 { \__fp_ep_inv_to_float_o:wwN 0, }

```

(End definition for __fp_fixed_inv_to_float_o:wwN.)

__fp_fixed_to_float_rad_o:wwN

Converts the fixed point number #1 from degrees to radians then to a floating point number. This could perhaps remain in l3fp-trig.

```

20187 \cs_new:Npn \__fp_fixed_to_float_rad_o:wwN #1;
20188 {
20189   \__fp_fixed_mul:wwn #1; {5729}{5779}{5130}{8232}{0876}{7981};
20190   { \__fp_ep_to_float_o:wwN 2, }
20191 }

```

(End definition for __fp_fixed_to_float_rad_o:wwN.)

```

\__fp_fixed_to_float_o:wN      ... \__fp_int_eval:w <exponent> \__fp_fixed_to_float_o:wN {\langle a_1 \rangle} {\langle a_2 \rangle} {\langle a_3 \rangle}
\__fp_fixed_to_float_o:Nw      {\langle a_4 \rangle} {\langle a_5 \rangle} {\langle a_6 \rangle} ; <sign>
                                yields
                                <exponent'> ; {\langle a'_1 \rangle} {\langle a'_2 \rangle} {\langle a'_3 \rangle} {\langle a'_4 \rangle} ;

```

And the `to_fixed` version gives six brace groups instead of 4, ensuring that $1000 \leq \langle a'_1 \rangle \leq 9999$. At this stage, we know that $\langle a_1 \rangle$ is positive (otherwise, it is sign of an error before), and we assume that it is less than 10^8 .¹⁰

```

20192 \cs_new:Npn \__fp_fixed_to_float_o:Nw #1#2;
20193 { \__fp_fixed_to_float_o:wN #2; #1 }
20194 \cs_new:Npn \__fp_fixed_to_float_o:wN #1#2#3#4#5#6; #7
20195 { % for the 8-digit-at-the-start thing
20196   + \__fp_int_eval:w \c__fp_block_int
20197   \exp_after:wN \exp_after:wN
20198   \exp_after:wN \__fp_fixed_to_loop:N
20199   \exp_after:wN \use_none:n
20200   \int_value:w \__fp_int_eval:w
20201   1 0000 0000 + #1 \exp_after:wN \__fp_use_none_stop_f:n
20202   \int_value:w 1#2 \exp_after:wN \__fp_use_none_stop_f:n
20203   \int_value:w 1#3#4 \exp_after:wN \__fp_use_none_stop_f:n
20204   \int_value:w 1#5#6
20205   \exp_after:wN ;
20206   \exp_after:wN ;
20207 }
20208 \cs_new:Npn \__fp_fixed_to_loop:N #1
20209 {
20210   \if_meaning:w 0 #1
20211   - 1
20212   \exp_after:wN \__fp_fixed_to_loop:N
20213   \else:
20214     \exp_after:wN \__fp_fixed_to_loop_end:w
20215     \exp_after:wN #1
20216   \fi:
20217 }
20218 \cs_new:Npn \__fp_fixed_to_loop_end:w #1 #2 ;
20219 {
20220   \if_meaning:w ; #1
20221     \exp_after:wN \__fp_fixed_to_float_zero:w
20222   \else:
20223     \exp_after:wN \__fp_pack_twice_four:wNNNNNNNN
20224     \exp_after:wN \__fp_pack_twice_four:wNNNNNNNN
20225     \exp_after:wN \__fp_fixed_to_float_pack:ww
20226     \exp_after:wN ;
20227   \fi:
20228   #1 #2 0000 0000 0000 0000 ;
20229 }
20230 \cs_new:Npn \__fp_fixed_to_float_zero:w ; 0000 0000 0000 0000 ;
20231 {
20232   - 2 * \c__fp_max_exponent_int ;
20233   {0000} {0000} {0000} {0000} ;
20234 }

```

¹⁰Bruno: I must double check this assumption.

```

20235 \cs_new:Npn \__fp_fixed_to_float_pack:ww #1 ; #2#3 ; ;
20236 {
20237   \if_int_compare:w #2 > 4 \exp_stop_f:
20238     \exp_after:wN \__fp_fixed_to_float_round_up:wnnnnw
20239   \fi:
20240   ; #1 ;
20241 }
20242 \cs_new:Npn \__fp_fixed_to_float_round_up:wnnnnw ; #1#2#3#4 ;
20243 {
20244   \exp_after:wN \__fp_basics_pack_high:NNNNNw
20245   \int_value:w \__fp_int_eval:w 1 #1#2
20246   \exp_after:wN \__fp_basics_pack_low:NNNNNw
20247   \int_value:w \__fp_int_eval:w 1 #3#4 + 1 ;
20248 }

```

(End definition for __fp_fixed_to_float_o:wN and __fp_fixed_to_float_o:Nw.)

```

20249 \</package>

```

33 l3fp-expo implementation

```

20250 \*package>
20251 \<@@=fp>

```

Unary functions.

```

\__fp_parse_word_exp:N
\__fp_parse_word_ln:N
\__fp_parse_word_fact:N
20252 \cs_new:Npn \__fp_parse_word_exp:N
20253 { \__fp_parse_unary_function:NNN \__fp_exp_o:w ? }
20254 \cs_new:Npn \__fp_parse_word_ln:N
20255 { \__fp_parse_unary_function:NNN \__fp_ln_o:w ? }
20256 \cs_new:Npn \__fp_parse_word_fact:N
20257 { \__fp_parse_unary_function:NNN \__fp_fact_o:w ? }

```

(End definition for __fp_parse_word_exp:N, __fp_parse_word_ln:N, and __fp_parse_word_fact:N.)

33.1 Logarithm

33.1.1 Work plan

As for many other functions, we filter out special cases in __fp_ln_o:w. Then __fp_ln_npos_o:w receives a positive normal number, which we write in the form $a \cdot 10^b$ with $a \in [0.1, 1)$.

The rest of this section is actually not in sync with the code. Or is the code not in sync with the section? In the current code, $c \in [1, 10]$ is such that $0.7 \leq ac < 1.4$.

We are given a positive normal number, of the form $a \cdot 10^b$ with $a \in [0.1, 1)$. To compute its logarithm, we find a small integer $5 \leq c < 50$ such that $0.91 \leq ac/5 < 1.1$, and use the relation

$$\ln(a \cdot 10^b) = b \cdot \ln(10) - \ln(c/5) + \ln(ac/5).$$

The logarithms $\ln(10)$ and $\ln(c/5)$ are looked up in a table. The last term is computed using the following Taylor series of \ln near 1:

$$\ln\left(\frac{ac}{5}\right) = \ln\left(\frac{1+t}{1-t}\right) = 2t \left(1 + t^2 \left(\frac{1}{3} + t^2 \left(\frac{1}{5} + t^2 \left(\frac{1}{7} + t^2 \left(\frac{1}{9} + \dots\right)\right)\right)\right)\right)$$

where $t = 1 - 10/(ac + 5)$. We can now see one reason for the choice of $ac \sim 5$: then $ac + 5 = 10(1 - \epsilon)$ with $-0.05 < \epsilon \leq 0.045$, hence

$$t = \frac{\epsilon}{1 - \epsilon} = \epsilon(1 + \epsilon)(1 + \epsilon^2)(1 + \epsilon^4) \dots,$$

is not too difficult to compute.

33.1.2 Some constants

A few values of the logarithm as extended fixed point numbers. Those are needed in the implementation. It turns out that we don't need the value of $\ln(5)$.

```
\c__fp_ln_i_fixed_t1
\c__fp_ln_ii_fixed_t1
\c__fp_ln_iii_fixed_t1
\c__fp_ln_iv_fixed_t1
\c__fp_ln_vi_fixed_t1
\c__fp_ln_vii_fixed_t1
\c__fp_ln_viii_fixed_t1
\c__fp_ln_ix_fixed_t1
\c__fp_ln_x_fixed_t1
20258 \tl_const:Nn \c__fp_ln_i_fixed_t1 { {0000}{0000}{0000}{0000}{0000}{0000};}
20259 \tl_const:Nn \c__fp_ln_ii_fixed_t1 { {6931}{4718}{0559}{9453}{0941}{7232};}
20260 \tl_const:Nn \c__fp_ln_iii_fixed_t1 { {10986}{1228}{8668}{1096}{9139}{5245};}
20261 \tl_const:Nn \c__fp_ln_iv_fixed_t1 { {13862}{9436}{1119}{8906}{1883}{4464};}
20262 \tl_const:Nn \c__fp_ln_vi_fixed_t1 { {17917}{5946}{9228}{0550}{0081}{2477};}
20263 \tl_const:Nn \c__fp_ln_vii_fixed_t1 { {19459}{1014}{9055}{3133}{0510}{5353};}
20264 \tl_const:Nn \c__fp_ln_viii_fixed_t1 { {20794}{4154}{1679}{8359}{2825}{1696};}
20265 \tl_const:Nn \c__fp_ln_ix_fixed_t1 { {21972}{2457}{7336}{2193}{8279}{0490};}
20266 \tl_const:Nn \c__fp_ln_x_fixed_t1 { {23025}{8509}{2994}{0456}{8401}{7991};}
```

(End definition for `\c__fp_ln_i_fixed_t1` and others.)

33.1.3 Sign, exponent, and special numbers

`__fp_ln_o:w` The logarithm of negative numbers (including $-\infty$ and -0) raises the “invalid” exception. The logarithm of $+0$ is $-\infty$, raising a division by zero exception. The logarithm of $+\infty$ or a `nan` is itself. Positive normal numbers call `__fp_ln_npos_o:w`.

```
20267 \cs_new:Npn \__fp_ln_o:w #1 \s__fp \__fp_chk:w #2#3#4; @
20268 {
20269   \if_meaning:w 2 #3
20270     \__fp_case_use:nw { \__fp_invalid_operation_o:nw { ln } }
20271   \fi:
20272   \if_case:w #2 \exp_stop_f:
20273     \__fp_case_use:nw
20274     { \__fp_division_by_zero_o:Nnw \c_minus_inf_fp { ln } }
20275   \or:
20276   \else:
20277     \__fp_case_return_same_o:w
20278   \fi:
20279   \__fp_ln_npos_o:w \s__fp \__fp_chk:w #2#3#4;
20280 }
```

(End definition for `__fp_ln_o:w`.)

33.1.4 Absolute ln

`__fp_ln_npos_o:w` We catch the case of a significand very close to 0.1 or to 1. In all other cases, the final result is at least 10^{-4} , and then an error of $0.5 \cdot 10^{-20}$ is acceptable.

```
20281 \cs_new:Npn \__fp_ln_npos_o:w \s__fp \__fp_chk:w 10#1#2#3;
20282 { %%^A todo: ln(1) should be "exact zero", not "underflow"
20283   \exp_after:wN \__fp_sanitize:Nw
20284   \int_value:w % for the overall sign
```

```

20285 \if_int_compare:w #1 < 1 \exp_stop_f:
20286 2
20287 \else:
20288 0
20289 \fi:
20290 \exp_after:wN \exp_stop_f:
20291 \int_value:w \__fp_int_eval:w % for the exponent
20292 \__fp_ln_significand:NNNNnnnnN #2#3
20293 \__fp_ln_exponent:wn {#1}
20294 }

```

(End definition for __fp_ln_npos_o:w.)

__fp_ln_significand:NNNNnnnnN __fp_ln_significand:NNNNnnnnN $\langle X_1 \rangle$ $\{\langle X_2 \rangle\}$ $\{\langle X_3 \rangle\}$ $\{\langle X_4 \rangle\}$ $\langle continuation \rangle$
This function expands to

$\langle continuation \rangle$ $\{\langle Y_1 \rangle\}$ $\{\langle Y_2 \rangle\}$ $\{\langle Y_3 \rangle\}$ $\{\langle Y_4 \rangle\}$ $\{\langle Y_5 \rangle\}$ $\{\langle Y_6 \rangle\}$;

where $Y = -\ln(X)$ as an extended fixed point.

```

20295 \cs_new:Npn \__fp_ln_significand:NNNNnnnnN #1#2#3#4
20296 {
20297 \exp_after:wN \__fp_ln_x_ii:wnnnnn
20298 \int_value:w
20299 \if_case:w #1 \exp_stop_f:
20300 \or:
20301 \if_int_compare:w #2 < 4 \exp_stop_f:
20302 \__fp_int_eval:w 10 - #2
20303 \else:
20304 6
20305 \fi:
20306 \or: 4
20307 \or: 3
20308 \or: 2
20309 \or: 2
20310 \or: 2
20311 \else: 1
20312 \fi:
20313 ; { #1 #2 #3 #4 }
20314 }

```

(End definition for __fp_ln_significand:NNNNnnnnN.)

__fp_ln_x_ii:wnnnnn We have thus found $c \in [1, 10]$ such that $0.7 \leq ac < 1.4$ in all cases. Compute $1 + x = 1 + ac \in [1.7, 2.4)$.

```

20315 \cs_new:Npn \__fp_ln_x_ii:wnnnnn #1; #2#3#4#5
20316 {
20317 \exp_after:wN \__fp_ln_div_after:Nw
20318 \cs:w c__fp_ln_ \__fp_int_to_roman:w #1 _fixed_tl \exp_after:wN \cs_end:
20319 \int_value:w
20320 \exp_after:wN \__fp_ln_x_iv:wnnnnnnnnn
20321 \int_value:w \__fp_int_eval:w
20322 \exp_after:wN \__fp_ln_x_iii_var:NNNNNw
20323 \int_value:w \__fp_int_eval:w 9999 9990 + #1*#2#3 +
20324 \exp_after:wN \__fp_ln_x_iii:NNNNNNw
20325 \int_value:w \__fp_int_eval:w 10 0000 0000 + #1*#4#5 ;

```

```

20326     {20000} {0000} {0000} {0000}
20327   } %^A todo: reoptimize (a generalization attempt failed).
20328 \cs_new:Npn \__fp_ln_x_iii:NNNNNw #1#2 #3#4#5#6 #7;
20329   { #1#2; {#3#4#5#6} {#7} }
20330 \cs_new:Npn \__fp_ln_x_iii_var:NNNNNw #1 #2#3#4#5 #6;
20331   {
20332     #1#2#3#4#5 + 1 ;
20333     {#1#2#3#4#5} {#6}
20334   }

```

The Taylor series to be used is expressed in terms of $t = (x - 1)/(x + 1) = 1 - 2/(x + 1)$. We now compute the quotient with extended precision, reusing some code from `__fp_/_o:ww`. Note that $1 + x$ is known exactly.

To reuse notations from `l3fp-basics`, we want to compute A/Z with $A = 2$ and $Z = x + 1$. In `l3fp-basics`, we considered the case where both A and Z are arbitrary, in the range $[0.1, 1)$, and we had to monitor the growth of the sequence of remainders A , B , C , etc. to ensure that no overflow occurred during the computation of the next quotient. The main source of risk was our choice to define the quotient as roughly $10^9 \cdot A/10^5 \cdot Z$: then A was bound to be below $2.147 \dots$, and this limit was never far.

In our case, we can simply work with $10^8 \cdot A$ and $10^4 \cdot Z$, because our reason to work with higher powers has gone: we needed the integer $y \simeq 10^5 \cdot Z$ to be at least 10^4 , and now, the definition $y \simeq 10^4 \cdot Z$ suffices.

Let us thus define $y = \lfloor 10^4 \cdot Z \rfloor + 1 \in (1.7 \cdot 10^4, 2.4 \cdot 10^4]$, and

$$Q_1 = \left\lfloor \frac{\lfloor 10^8 \cdot A \rfloor}{y} - \frac{1}{2} \right\rfloor.$$

(The $1/2$ comes from how ε -TeX rounds.) As for division, it is easy to see that $Q_1 \leq 10^4 A/Z$, *i.e.*, Q_1 is an underestimate.

Exactly as we did for division, we set $B = 10^4 A - Q_1 Z$. Then

$$\begin{aligned}
10^4 B &\leq A_1 A_2 \cdot A_3 A_4 - \left(\frac{A_1 A_2}{y} - \frac{3}{2} \right) 10^4 Z \\
&\leq A_1 A_2 \left(1 - \frac{10^4 Z}{y} \right) + 1 + \frac{3}{2} y \\
&\leq 10^8 \frac{A}{y} + 1 + \frac{3}{2} y
\end{aligned}$$

In the same way, and using $1.7 \cdot 10^4 \leq y \leq 2.4 \cdot 10^4$, and convexity, we get

$$\begin{aligned} 10^4 A &= 2 \cdot 10^4 \\ 10^4 B &\leq 10^8 \frac{A}{y} + 1.6y \leq 4.7 \cdot 10^4 \\ 10^4 C &\leq 10^8 \frac{B}{y} + 1.6y \leq 5.8 \cdot 10^4 \\ 10^4 D &\leq 10^8 \frac{C}{y} + 1.6y \leq 6.3 \cdot 10^4 \\ 10^4 E &\leq 10^8 \frac{D}{y} + 1.6y \leq 6.5 \cdot 10^4 \\ 10^4 F &\leq 10^8 \frac{E}{y} + 1.6y \leq 6.6 \cdot 10^4 \end{aligned}$$

Note that we compute more steps than for division: since t is not the end result, we need to know it with more accuracy (on the other hand, the ending is much simpler, as we don't need an exact rounding for transcendental functions, but just a faithful rounding).

`__fp_ln_x_iv:wnnnnnnnn <1 or 2> <8d> ; {<4d>} {<4d>} <fixed-t1>`

The number is x . Compute y by adding 1 to the five first digits.

```

20335 \cs_new:Npn __fp_ln_x_iv:wnnnnnnnn #1; #2#3#4#5 #6#7#8#9
20336 {
20337   \exp_after:wN __fp_div_significand_pack:NNN
20338   \int_value:w __fp_int_eval:w
20339   __fp_ln_div_i:w #1 ;
20340   #6 #7 ; {#8} {#9}
20341   {#2} {#3} {#4} {#5}
20342   { \exp_after:wN __fp_ln_div_ii:wnn \int_value:w #1 }
20343   { \exp_after:wN __fp_ln_div_ii:wnn \int_value:w #1 }
20344   { \exp_after:wN __fp_ln_div_ii:wnn \int_value:w #1 }
20345   { \exp_after:wN __fp_ln_div_ii:wnn \int_value:w #1 }
20346   { \exp_after:wN __fp_ln_div_vi:wnn \int_value:w #1 }
20347 }
20348 \cs_new:Npn __fp_ln_div_i:w #1;
20349 {
20350   \exp_after:wN __fp_div_significand_calc:wnnnnnnnn
20351   \int_value:w __fp_int_eval:w 999999 + 2 0000 0000 / #1 ; % Q1
20352 }
20353 \cs_new:Npn __fp_ln_div_ii:wnn #1; #2;#3 % y; B1;B2 <- for k=1
20354 {
20355   \exp_after:wN __fp_div_significand_pack:NNN
20356   \int_value:w __fp_int_eval:w
20357   \exp_after:wN __fp_div_significand_calc:wnnnnnnnn
20358   \int_value:w __fp_int_eval:w 999999 + #2 #3 / #1 ; % Q2
20359   #2 #3 ;
20360 }
20361 \cs_new:Npn __fp_ln_div_vi:wnn #1; #2;#3#4#5 #6#7#8#9 %y;F1;F2F3F4x1x2x3x4
20362 {
20363   \exp_after:wN __fp_div_significand_pack:NNN

```

```

20364 \int_value:w \_fp_int_eval:w 1000000 + #2 #3 / #1 ; % Q6
20365 }

```

We now have essentially

```

\_fp_ln_div_after:Nw <fixed t1>
\_fp_div_significand_pack:NNN 106 + Q1
\_fp_div_significand_pack:NNN 106 + Q2
\_fp_div_significand_pack:NNN 106 + Q3
\_fp_div_significand_pack:NNN 106 + Q4
\_fp_div_significand_pack:NNN 106 + Q5
\_fp_div_significand_pack:NNN 106 + Q6 ;
<exponent> ; <continuation>

```

where $\langle fixed\ t1 \rangle$ holds the logarithm of a number in $[1, 10]$, and $\langle exponent \rangle$ is the exponent. Also, the expansion is done backwards. Then `_fp_div_significand_pack:NNN` puts things in the correct order to add the Q_i together and put semicolons between each piece. Once those have been expanded, we get

```

\_fp_ln_div_after:Nw <fixed-t1> <1d> ; <4d> ; <4d> ;
<4d> ; <4d> ; <4d> ; <4d> ; <exponent> ;

```

Just as with division, we know that the first two digits are 1 and 0 because of bounds on the final result of the division $2/(x+1)$, which is between roughly 0.8 and 1.2. We then compute $1 - 2/(x+1)$, after testing whether $2/(x+1)$ is greater than or smaller than 1.

```

20366 \cs_new:Npn \_fp_ln_div_after:Nw #1#2;
20367 {
20368   \if_meaning:w 0 #2
20369   \exp_after:wN \_fp_ln_t_small:Nw
20370   \else:
20371   \exp_after:wN \_fp_ln_t_large:NNw
20372   \exp_after:wN -
20373   \fi:
20374   #1
20375 }
20376 \cs_new:Npn \_fp_ln_t_small:Nw #1 #2; #3; #4; #5; #6; #7;
20377 {
20378   \exp_after:wN \_fp_ln_t_large:NNw
20379   \exp_after:wN + % <sign>
20380   \exp_after:wN #1
20381   \int_value:w \_fp_int_eval:w 9999 - #2 \exp_after:wN ;
20382   \int_value:w \_fp_int_eval:w 9999 - #3 \exp_after:wN ;
20383   \int_value:w \_fp_int_eval:w 9999 - #4 \exp_after:wN ;
20384   \int_value:w \_fp_int_eval:w 9999 - #5 \exp_after:wN ;
20385   \int_value:w \_fp_int_eval:w 9999 - #6 \exp_after:wN ;
20386   \int_value:w \_fp_int_eval:w 1 0000 - #7 ;
20387 }

\_fp_ln_t_large:NNw <sign> <fixed t1>
<t123456

```

Compute the square t^2 , and keep t at the end with its sign. We know that $t < 0.1765$, so every piece has at most 4 digits. However, since we were not careful in `_fp_ln_t_small:w`, they can have less than 4 digits.

```

20388 \cs_new:Npn \__fp_ln_t_large:NNw #1 #2 #3; #4; #5; #6; #7; #8;
20389 {
20390   \exp_after:wN \__fp_ln_square_t_after:w
20391   \int_value:w \__fp_int_eval:w 9999 0000 + #3*#3
20392   \exp_after:wN \__fp_ln_square_t_pack:NNNNNw
20393   \int_value:w \__fp_int_eval:w 9999 0000 + 2*#3*#4
20394   \exp_after:wN \__fp_ln_square_t_pack:NNNNNw
20395   \int_value:w \__fp_int_eval:w 9999 0000 + 2*#3*#5 + #4*#4
20396   \exp_after:wN \__fp_ln_square_t_pack:NNNNNw
20397   \int_value:w \__fp_int_eval:w 9999 0000 + 2*#3*#6 + 2*#4*#5
20398   \exp_after:wN \__fp_ln_square_t_pack:NNNNNw
20399   \int_value:w \__fp_int_eval:w
20400     1 0000 0000 + 2*#3*#7 + 2*#4*#6 + #5*#5
20401     + (2*#3*#8 + 2*#4*#7 + 2*#5*#6) / 1 0000
20402     % ; ; ;
20403   \exp_after:wN \__fp_ln_twice_t_after:w
20404   \int_value:w \__fp_int_eval:w -1 + 2*#3
20405   \exp_after:wN \__fp_ln_twice_t_pack:Nw
20406   \int_value:w \__fp_int_eval:w 9999 + 2*#4
20407   \exp_after:wN \__fp_ln_twice_t_pack:Nw
20408   \int_value:w \__fp_int_eval:w 9999 + 2*#5
20409   \exp_after:wN \__fp_ln_twice_t_pack:Nw
20410   \int_value:w \__fp_int_eval:w 9999 + 2*#6
20411   \exp_after:wN \__fp_ln_twice_t_pack:Nw
20412   \int_value:w \__fp_int_eval:w 9999 + 2*#7
20413   \exp_after:wN \__fp_ln_twice_t_pack:Nw
20414   \int_value:w \__fp_int_eval:w 10000 + 2*#8 ; ;
20415   { \__fp_ln_c:NwNw #1 }
20416   #2
20417 }
20418 \cs_new:Npn \__fp_ln_twice_t_pack:Nw #1 #2; { + #1 ; {#2} }
20419 \cs_new:Npn \__fp_ln_twice_t_after:w #1; { ; ; ; {#1} }
20420 \cs_new:Npn \__fp_ln_square_t_pack:NNNNNw #1 #2#3#4#5 #6;
20421   { + #1#2#3#4#5 ; {#6} }
20422 \cs_new:Npn \__fp_ln_square_t_after:w 1 0 #1#2#3 #4;
20423   { \__fp_ln_Taylor:wwNw {0#1#2#3} {#4} }

```

(End definition for __fp_ln_x_ii:wnnnn.)

__fp_ln_Taylor:wwNw Denoting $T = t^2$, we get

```

\__fp_ln_Taylor:wwNw
{ \langle T_1 \rangle } { \langle T_2 \rangle } { \langle T_3 \rangle } { \langle T_4 \rangle } { \langle T_5 \rangle } { \langle T_6 \rangle } ; ;
{ \langle (2t)_1 \rangle } { \langle (2t)_2 \rangle } { \langle (2t)_3 \rangle } { \langle (2t)_4 \rangle } { \langle (2t)_5 \rangle } { \langle (2t)_6 \rangle } ;
{ \__fp_ln_c:NwNw \langle sign \rangle }
\langle fixed t1 \rangle \langle exponent \rangle ; \langle continuation \rangle

```

And we want to compute

$$\ln\left(\frac{1+t}{1-t}\right) = 2t \left(1 + T \left(\frac{1}{3} + T \left(\frac{1}{5} + T \left(\frac{1}{7} + T \left(\frac{1}{9} + \cdots \right) \right) \right) \right) \right)$$

The process looks as follows

```

\loop 5; A;
\div_int 5; 1.0; \add A; \mul T; {\loop \eval 5-2;}
\add 0.2; A; \mul T; {\loop \eval 5-2;}
\mul B; T; {\loop 3;}
\loop 3; C;

```

This uses the routine for dividing a number by a small integer ($< 10^4$).

```

20424 \cs_new:Npn \__fp_ln_Taylor:wwNw
20425   { \__fp_ln_Taylor_loop:www 21 ; {0000}{0000}{0000}{0000}{0000}{0000} ; }
20426 \cs_new:Npn \__fp_ln_Taylor_loop:www #1; #2; #3;
20427   {
20428     \if_int_compare:w #1 = 1 \exp_stop_f:
20429     \__fp_ln_Taylor_break:w
20430     \fi:
20431     \exp_after:wN \__fp_fixed_div_int:wwN \c__fp_one_fixed_tl #1;
20432     \__fp_fixed_add:wwn #2;
20433     \__fp_fixed_mul:wwn #3;
20434     {
20435       \exp_after:wN \__fp_ln_Taylor_loop:www
20436       \int_value:w \__fp_int_eval:w #1 - 2 ;
20437     }
20438     #3;
20439   }
20440 \cs_new:Npn \__fp_ln_Taylor_break:w \fi: #1 \__fp_fixed_add:wwn #2#3; #4 ;;
20441   {
20442     \fi:
20443     \exp_after:wN \__fp_fixed_mul:wwn
20444     \exp_after:wN { \int_value:w \__fp_int_eval:w 10000 + #2 } #3;
20445   }

```

(End definition for $\backslash_fp_ln_Taylor:wwNw$.)

```

\__fp_ln_c:NwNw <sign>
{\langle r_1 \rangle} {\langle r_2 \rangle} {\langle r_3 \rangle} {\langle r_4 \rangle} {\langle r_5 \rangle} {\langle r_6 \rangle} ;
<fixed tl> <exponent> ; <continuation>

```

We are now reduced to finding $\ln(c)$ and $\langle exponent \rangle \ln(10)$ in a table, and adding it to the mixture. The first step is to get $\ln(c) - \ln(x) = -\ln(a)$, then we get $\mathbf{b} \ln(10)$ and add or subtract.

For now, $\ln(x)$ is given as $\cdot 10^0$. Unless both the exponent is 1 and $c = 1$, we shift to working in units of $\cdot 10^4$, since the final result is at least $\ln(10/7) \simeq 0.35$.

```

20446 \cs_new:Npn \__fp_ln_c:NwNw #1 #2; #3
20447   {
20448     \if_meaning:w + #1
20449     \exp_after:wN \exp_after:wN \exp_after:wN \__fp_fixed_sub:wwn
20450     \else:
20451     \exp_after:wN \exp_after:wN \exp_after:wN \__fp_fixed_add:wwn
20452     \fi:
20453     #3 #2 ;
20454   }

```

(End definition for $\backslash_fp_ln_c:NwNw$.)

```

    \_fp_ln_exponent:wn
    {\langle s_1 \rangle} {\langle s_2 \rangle} {\langle s_3 \rangle} {\langle s_4 \rangle} {\langle s_5 \rangle} {\langle s_6 \rangle} ;
    {\langle exponent \rangle}

```

Compute $\langle exponent \rangle$ times $\ln(10)$. Apart from the cases where $\langle exponent \rangle$ is 0 or 1, the result is necessarily at least $\ln(10) \simeq 2.3$ in magnitude. We can thus drop the least significant 4 digits. In the case of a very large (positive or negative) exponent, we can (and we need to) drop 4 additional digits, since the result is of order 10^4 . Naively, one would think that in both cases we can drop 4 more digits than we do, but that would be slightly too tight for rounding to happen correctly. Besides, we already have addition and subtraction for 24 digits fixed point numbers.

```

20455 \cs_new:Npn \_fp_ln_exponent:wn #1; #2
20456 {
20457   \if_case:w #2 \exp_stop_f:
20458     0 \_fp_case_return:nw { \_fp_fixed_to_float_o:Nw 2 }
20459   \or:
20460     \exp_after:wN \_fp_ln_exponent_one:ww \int_value:w
20461   \else:
20462     \if_int_compare:w #2 > 0 \exp_stop_f:
20463     \exp_after:wN \_fp_ln_exponent_small:NNww
20464     \exp_after:wN 0
20465     \exp_after:wN \_fp_fixed_sub:wwn \int_value:w
20466   \else:
20467     \exp_after:wN \_fp_ln_exponent_small:NNww
20468     \exp_after:wN 2
20469     \exp_after:wN \_fp_fixed_add:wwn \int_value:w -
20470   \fi:
20471   \fi:
20472   #2; #1;
20473 }

```

Now we painfully write all the cases.¹¹ No overflow nor underflow can happen, except when computing $\ln(1)$.

```

20474 \cs_new:Npn \_fp_ln_exponent_one:ww 1; #1;
20475 {
20476   0
20477   \exp_after:wN \_fp_fixed_sub:wwn \c__fp_ln_x_fixed_tl #1;
20478   \_fp_fixed_to_float_o:wN 0
20479 }

```

For small exponents, we just drop one block of digits, and set the exponent of the log to 4 (minus any shift coming from leading zeros in the conversion from fixed point to floating point). Note that here the exponent has been made positive.

```

20480 \cs_new:Npn \_fp_ln_exponent_small:NNww #1#2#3; #4#5#6#7#8#9;
20481 {
20482   4
20483   \exp_after:wN \_fp_fixed_mul:wwn
20484     \c__fp_ln_x_fixed_tl
20485     {#3}{0000}{0000}{0000}{0000}{0000} ;
20486   #2
20487   {0000}{#4}{#5}{#6}{#7}{#8};
20488   \_fp_fixed_to_float_o:wN #1
20489 }

```

¹¹Bruno: do rounding.

(End definition for _fp_ln_exponent:wn.)

33.2 Exponential

33.2.1 Sign, exponent, and special numbers

_fp_exp_o:w

```

20490 \cs_new:Npn \_fp_exp_o:w #1 \s__fp \_fp_chk:w #2#3#4; @
20491 {
20492   \if_case:w #2 \exp_stop_f:
20493     \_fp_case_return_o:Nw \c_one_fp
20494   \or:
20495     \exp_after:wN \_fp_exp_normal_o:w
20496   \or:
20497     \if_meaning:w 0 #3
20498       \exp_after:wN \_fp_case_return_o:Nw
20499       \exp_after:wN \c_inf_fp
20500     \else:
20501       \exp_after:wN \_fp_case_return_o:Nw
20502       \exp_after:wN \c_zero_fp
20503     \fi:
20504   \or:
20505     \_fp_case_return_same_o:w
20506   \fi:
20507   \s__fp \_fp_chk:w #2#3#4;
20508 }

```

(End definition for _fp_exp_o:w.)

_fp_exp_normal_o:w

_fp_exp_pos_o:NNwnw

_fp_exp_overflow:NN

```

20509 \cs_new:Npn \_fp_exp_normal_o:w \s__fp \_fp_chk:w 1#1
20510 {
20511   \if_meaning:w 0 #1
20512     \_fp_exp_pos_o:NNwnw + \_fp_fixed_to_float_o:wN
20513   \else:
20514     \_fp_exp_pos_o:NNwnw - \_fp_fixed_inv_to_float_o:wN
20515   \fi:
20516 }
20517 \cs_new:Npn \_fp_exp_pos_o:NNwnw #1#2#3 \fi: #4#5;
20518 {
20519   \fi:
20520   \if_int_compare:w #4 > \c__fp_max_exp_exponent_int
20521     \token_if_eq_charcode:NNTF + #1
20522     { \_fp_exp_overflow:NN \_fp_overflow:w \c_inf_fp }
20523     { \_fp_exp_overflow:NN \_fp_underflow:w \c_zero_fp }
20524   \exp:w
20525   \else:
20526     \exp_after:wN \_fp_sanitize:Nw
20527     \exp_after:wN 0
20528     \int_value:w #1 \_fp_int_eval:w
20529     \if_int_compare:w #4 < 0 \exp_stop_f:
20530       \exp_after:wN \use_i:nn
20531     \else:
20532       \exp_after:wN \use_ii:nn

```

```

20533     \fi:
20534     {
20535         0
20536         \__fp_decimate:nNnnnn { - #4 }
20537         \__fp_exp_Taylor:Nnnwn
20538     }
20539     {
20540         \__fp_decimate:nNnnnn { \c__fp_prec_int - #4 }
20541         \__fp_exp_pos_large:NnnNwn
20542     }
20543     #5
20544     {#4}
20545     #1 #2 0
20546     \exp:w
20547     \fi:
20548     \exp_after:wN \exp_end:
20549 }
20550 \cs_new:Npn \__fp_exp_overflow:NN #1#2
20551 {
20552     \exp_after:wN \exp_after:wN
20553     \exp_after:wN #1
20554     \exp_after:wN #2
20555 }

```

(End definition for __fp_exp_normal_o:w, __fp_exp_pos_o:Nnnwn, and __fp_exp_overflow:NN.)

```

\__fp_exp_Taylor:Nnnwn
\__fp_exp_Taylor_loop:www
\__fp_exp_Taylor_break:Nww

```

This function is called for numbers in the range $[10^{-9}, 10^{-1}]$. We compute 10 terms of the Taylor series. The first argument is irrelevant (rounding digit used by some other functions). The next three arguments, at least 16 digits, delimited by a semicolon, form a fixed point number, so we pack it in blocks of 4 digits.

```

20556 \cs_new:Npn \__fp_exp_Taylor:Nnnwn #1#2#3 #4; #5 #6
20557 {
20558     #6
20559     \__fp_pack_twice_four:wNNNNNNNN
20560     \__fp_pack_twice_four:wNNNNNNNN
20561     \__fp_pack_twice_four:wNNNNNNNN
20562     \__fp_exp_Taylor_ii:ww
20563     ; #2#3#4 0000 0000 ;
20564 }
20565 \cs_new:Npn \__fp_exp_Taylor_ii:ww #1; #2;
20566 { \__fp_exp_Taylor_loop:www 10 ; #1 ; #1 ; \s__fp_stop }
20567 \cs_new:Npn \__fp_exp_Taylor_loop:www #1; #2; #3;
20568 {
20569     \if_int_compare:w #1 = 1 \exp_stop_f:
20570     \exp_after:wN \__fp_exp_Taylor_break:Nww
20571     \fi:
20572     \__fp_fixed_div_int:wwN #3 ; #1 ;
20573     \__fp_fixed_add_one:wN
20574     \__fp_fixed_mul:wwN #2 ;
20575     {
20576         \exp_after:wN \__fp_exp_Taylor_loop:www
20577         \int_value:w \__fp_int_eval:w #1 - 1 ;
20578         #2 ;
20579     }

```

```

20580     }
20581 \cs_new:Npn \__fp_exp_Taylor_break:Nww #1 #2; #3 \s__fp_stop
20582 { \__fp_fixed_add_one:wN #2 ; }

```

(End definition for __fp_exp_Taylor:Nnnwn, __fp_exp_Taylor_loop:www, and __fp_exp_Taylor-break:Nww.)

\c__fp_exp_intarray The integer array has $6 \times 9 \times 4 = 216$ items encoding the values of $\exp(j \times 10^i)$ for $j = 1, \dots, 9$ and $i = -1, \dots, 4$. Each value is expressed as $\simeq 10^p \times 0.m_1m_2m_3$ with three 8-digit blocks m_1, m_2, m_3 and an integer exponent p (one more than the scientific exponent), and these are stored in the integer array as four items: $p, 10^8 + m_1, 10^8 + m_2, 10^8 + m_3$. The various exponentials are stored in increasing order of $j \times 10^i$.

Storing this data in an integer array makes it slightly harder to access (slower, too), but uses 16 bytes of memory per exponential stored, while storing as tokens used around 40 tokens; tokens have an especially large footprint in Unicode-aware engines.

```

20583 \intarray_const_from_clist:Nn \c__fp_exp_intarray
20584 {
20585     1 , 1 1105 1709 , 1 1807 5647 , 1 6248 1171 ,
20586     1 , 1 1221 4027 , 1 5816 0169 , 1 8339 2107 ,
20587     1 , 1 1349 8588 , 1 0757 6003 , 1 1039 8374 ,
20588     1 , 1 1491 8246 , 1 9764 1270 , 1 3178 2485 ,
20589     1 , 1 1648 7212 , 1 7070 0128 , 1 1468 4865 ,
20590     1 , 1 1822 1188 , 1 0039 0508 , 1 9748 7537 ,
20591     1 , 1 2013 7527 , 1 0747 0476 , 1 5216 2455 ,
20592     1 , 1 2225 5409 , 1 2849 2467 , 1 6045 7954 ,
20593     1 , 1 2459 6031 , 1 1115 6949 , 1 6638 0013 ,
20594     1 , 1 2718 2818 , 1 2845 9045 , 1 2353 6029 ,
20595     1 , 1 7389 0560 , 1 9893 0650 , 1 2272 3043 ,
20596     2 , 1 2008 5536 , 1 9231 8766 , 1 7740 9285 ,
20597     2 , 1 5459 8150 , 1 0331 4423 , 1 9078 1103 ,
20598     3 , 1 1484 1315 , 1 9102 5766 , 1 0342 1116 ,
20599     3 , 1 4034 2879 , 1 3492 7351 , 1 2260 8387 ,
20600     4 , 1 1096 6331 , 1 5842 8458 , 1 5992 6372 ,
20601     4 , 1 2980 9579 , 1 8704 1728 , 1 2747 4359 ,
20602     4 , 1 8103 0839 , 1 2757 5384 , 1 0077 1000 ,
20603     5 , 1 2202 6465 , 1 7948 0671 , 1 6516 9579 ,
20604     9 , 1 4851 6519 , 1 5409 7902 , 1 7796 9107 ,
20605     14 , 1 1068 6474 , 1 5815 2446 , 1 2146 9905 ,
20606     18 , 1 2353 8526 , 1 6837 0199 , 1 8540 7900 ,
20607     22 , 1 5184 7055 , 1 2858 7072 , 1 4640 8745 ,
20608     27 , 1 1142 0073 , 1 8981 5684 , 1 2836 6296 ,
20609     31 , 1 2515 4386 , 1 7091 9167 , 1 0062 6578 ,
20610     35 , 1 5540 6223 , 1 8439 3510 , 1 0525 7117 ,
20611     40 , 1 1220 4032 , 1 9431 7840 , 1 8020 0271 ,
20612     44 , 1 2688 1171 , 1 4181 6135 , 1 4484 1263 ,
20613     87 , 1 7225 9737 , 1 6812 5749 , 1 2581 7748 ,
20614     131 , 1 1942 4263 , 1 9524 1255 , 1 9365 8421 ,
20615     174 , 1 5221 4696 , 1 8976 4143 , 1 9505 8876 ,
20616     218 , 1 1403 5922 , 1 1785 2837 , 1 4107 3977 ,
20617     261 , 1 3773 0203 , 1 0092 9939 , 1 8234 0143 ,
20618     305 , 1 1014 2320 , 1 5473 5004 , 1 5094 5533 ,
20619     348 , 1 2726 3745 , 1 7211 2566 , 1 5673 6478 ,
20620     391 , 1 7328 8142 , 1 2230 7421 , 1 7051 8866 ,
20621     435 , 1 1970 0711 , 1 1401 7046 , 1 9938 8888 ,

```

```

20622      869 , 1 3881 1801 , 1 9428 4368 , 1 5764 8232 ,
20623      1303 , 1 7646 2009 , 1 8905 4704 , 1 8893 1073 ,
20624      1738 , 1 1506 3559 , 1 7005 0524 , 1 9009 7592 ,
20625      2172 , 1 2967 6283 , 1 8402 3667 , 1 0689 6630 ,
20626      2606 , 1 5846 4389 , 1 5650 2114 , 1 7278 5046 ,
20627      3041 , 1 1151 7900 , 1 5080 6878 , 1 2914 4154 ,
20628      3475 , 1 2269 1083 , 1 0850 6857 , 1 8724 4002 ,
20629      3909 , 1 4470 3047 , 1 3316 5442 , 1 6408 6591 ,
20630      4343 , 1 8806 8182 , 1 2566 2921 , 1 5872 6150 ,
20631      8686 , 1 7756 0047 , 1 2598 6861 , 1 0458 3204 ,
20632      13029 , 1 6830 5723 , 1 7791 4884 , 1 1932 7351 ,
20633      17372 , 1 6015 5609 , 1 3095 3052 , 1 3494 7574 ,
20634      21715 , 1 5297 7951 , 1 6443 0315 , 1 3251 3576 ,
20635      26058 , 1 4665 6719 , 1 0099 3379 , 1 5527 2929 ,
20636      30401 , 1 4108 9724 , 1 3326 3186 , 1 5271 5665 ,
20637      34744 , 1 3618 6973 , 1 3140 0875 , 1 3856 4102 ,
20638      39087 , 1 3186 9209 , 1 6113 3900 , 1 6705 9685 ,
20639    }

```

(End definition for \c__fp_exp_intarray.)

__fp_exp_pos_large:NnnNwn The first two arguments are irrelevant (a rounding digit, and a brace group with 8 zeros).
 __fp_exp_large_after:wnn The third argument is the integer part of our number, then we have the decimal part
 __fp_exp_large:NwN delimited by a semicolon, and finally the exponent, in the range [0,5]. Remove leading
 __fp_exp_intarray:w zeros from the integer part: putting #4 in there too ensures that an integer part of 0 is
 __fp_exp_intarray_aux:w also removed. Then read digits one by one, looking up $\exp(\langle digit \rangle \cdot 10^{\langle exponent \rangle})$ in a table,
 and multiplying that to the current total. The loop is done by __fp_exp_large:NwN,
 whose #1 is the $\langle exponent \rangle$, #2 is the current mantissa, and #3 is the $\langle digit \rangle$. At the end,
 __fp_exp_large_after:wnn moves on to the Taylor series, eventually multiplied with
 the mantissa that we have just computed.

```

20640 \cs_new:Npn \__fp_exp_pos_large:NnnNwn #1#2#3 #4#5; #6
20641 {
20642   \exp_after:wN \exp_after:wN \exp_after:wN \__fp_exp_large:NwN
20643   \exp_after:wN \exp_after:wN \exp_after:wN #6
20644   \exp_after:wN \c__fp_one_fixed_tl
20645   \int_value:w #3 #4 \exp_stop_f:
20646   #5 00000 ;
20647 }
20648 \cs_new:Npn \__fp_exp_large:NwN #1#2; #3
20649 {
20650   \if_case:w #3 ~
20651     \exp_after:wN \__fp_fixed_continue:wn
20652   \else:
20653     \exp_after:wN \__fp_exp_intarray:w
20654     \int_value:w \__fp_int_eval:w 36 * #1 + 4 * #3 \exp_after:wN ;
20655   \fi:
20656   #2;
20657   {
20658     \if_meaning:w 0 #1
20659     \exp_after:wN \__fp_exp_large_after:wnn
20660   \else:
20661     \exp_after:wN \__fp_exp_large:NwN
20662     \int_value:w \__fp_int_eval:w #1 - 1 \exp_after:wN \scan_stop:
20663   \fi:

```

```

20664     }
20665   }
20666   \cs_new:Npn \__fp_exp_intarray:w #1 ;
20667   {
20668     +
20669     \__kernel_intarray_item:Nn \c__fp_exp_intarray
20670     { \__fp_int_eval:w #1 - 3 \scan_stop: }
20671     \exp_after:wN \use_i:nnn
20672     \exp_after:wN \__fp_fixed_mul:wwn
20673     \int_value:w 0
20674     \exp_after:wN \__fp_exp_intarray_aux:w
20675     \int_value:w \__kernel_intarray_item:Nn
20676     \c__fp_exp_intarray { \__fp_int_eval:w #1 - 2 }
20677     \exp_after:wN \__fp_exp_intarray_aux:w
20678     \int_value:w \__kernel_intarray_item:Nn
20679     \c__fp_exp_intarray { \__fp_int_eval:w #1 - 1 }
20680     \exp_after:wN \__fp_exp_intarray_aux:w
20681     \int_value:w \__kernel_intarray_item:Nn \c__fp_exp_intarray {#1} ; ;
20682   }
20683   \cs_new:Npn \__fp_exp_intarray_aux:w 1 #1#2#3#4#5 ; { ; {#1#2#3#4} {#5} }
20684   \cs_new:Npn \__fp_exp_large_after:wwn #1; #2; #3
20685   {
20686     \__fp_exp_Taylor:Nnnwn ? { } { } 0 #2; { } #3
20687     \__fp_fixed_mul:wwn #1;
20688   }

```

(End definition for `__fp_exp_pos_large:NnnNwn` and others.)

33.3 Power

Raising a number a to a power b leads to many distinct situations.

a^b	$-\infty$	$(-\infty, -0)$	$-\text{integer}$	± 0	$+\text{integer}$	$(0, \infty)$	$+\infty$	NaN
$+\infty$	$+0$		$+0$	$+1$	$+\infty$		$+\infty$	NaN
$(1, \infty)$	$+0$		$+ a ^b$	$+1$	$+ a ^b$		$+\infty$	NaN
$+1$	$+1$		$+1$	$+1$	$+1$		$+1$	$+1$
$(0, 1)$	$+\infty$		$+ a ^b$	$+1$	$+ a ^b$		$+0$	NaN
$+0$	$+\infty$		$+\infty$	$+1$	$+0$		$+0$	NaN
-0	$+\infty$	NaN	$(-1)^b \infty$	$+1$	$(-1)^b 0$	$+0$	$+0$	NaN
$(-1, 0)$	$+\infty$	NaN	$(-1)^b a ^b$	$+1$	$(-1)^b a ^b$	NaN	$+0$	NaN
-1	$+1$	NaN	$(-1)^b$	$+1$	$(-1)^b$	NaN	$+1$	NaN
$(-\infty, -1)$	$+0$	NaN	$(-1)^b a ^b$	$+1$	$(-1)^b a ^b$	NaN	$+\infty$	NaN
$-\infty$	$+0$	$+0$	$(-1)^b 0$	$+1$	$(-1)^b \infty$	NaN	$+\infty$	NaN
NaN	NaN	NaN	NaN	$+1$	NaN	NaN	NaN	NaN

We distinguished in this table the cases of finite (positive or negative) integer exponents, as $(-1)^b$ is defined in that case. One peculiarity of this operation is that $\text{NaN}^0 = 1^{\text{NaN}} = 1$, because this relation is obeyed for any number, even $\pm\infty$.

`__fp_~_o:ww` We cram most of the tests into a single function to save csnames. First treat the case $b = 0$: $a^0 = 1$ for any a , even `nan`. Then test the sign of a .

- If it is positive, and a is a normal number, call `__fp_pow_normal_o:ww` followed by the two `fp` a and b . For $a = +0$ or $+\infty$, call `__fp_pow_zero_or_inf:ww` instead, to return either $+0$ or $+\infty$ as appropriate.
- If a is a `nan`, then skip to the next semicolon (which happens to be conveniently the end of b) and return `nan`.
- Finally, if a is negative, compute a^b (`__fp_pow_normal_o:ww` which ignores the sign of its first operand), and keep an extra copy of a and b (the second brace group, containing $\{ b a \}$, is inserted between a and b). Then do some tests to find the final sign of the result if it exists.

```

20689 \cs_new:cpn { __fp_ \iow_char:N \^_ _o:ww }
20690   \s__fp \__fp_chk:w #1#2#3; \s__fp \__fp_chk:w #4#5#6;
20691   {
20692     \if_meaning:w 0 #4
20693       \__fp_case_return_o:Nw \c_one_fp
20694     \fi:
20695     \if_case:w #2 \exp_stop_f:
20696       \exp_after:wN \use_i:nn
20697     \or:
20698       \__fp_case_return_o:Nw \c_nan_fp
20699     \else:
20700       \exp_after:wN \__fp_pow_neg:www
20701       \exp:w \exp_end_continue_f:w \exp_after:wN \use:nn
20702     \fi:
20703     {
20704       \if_meaning:w 1 #1
20705         \exp_after:wN \__fp_pow_normal_o:ww
20706       \else:
20707         \exp_after:wN \__fp_pow_zero_or_inf:ww
20708       \fi:
20709       \s__fp \__fp_chk:w #1#2#3;
20710     }
20711     { \s__fp \__fp_chk:w #4#5#6; \s__fp \__fp_chk:w #1#2#3; }
20712     \s__fp \__fp_chk:w #4#5#6;
20713   }

```

(End definition for `__fp_~_o:ww`.)

`__fp_pow_zero_or_inf:ww` Raising -0 or $-\infty$ to `nan` yields `nan`. For other powers, the result is $+0$ if 0 is raised to a positive power or ∞ to a negative power, and $+\infty$ otherwise. Thus, if the type of a and the sign of b coincide, the result is 0 , since those conveniently take the same possible values, 0 and 2 . Otherwise, either $a = \pm\infty$ and $b > 0$ and the result is $+\infty$, or $a = \pm 0$ with $b < 0$ and we have a division by zero unless $b = -\infty$.

```

20714 \cs_new:Npn \__fp_pow_zero_or_inf:ww
20715   \s__fp \__fp_chk:w #1#2; \s__fp \__fp_chk:w #3#4
20716   {
20717     \if_meaning:w 1 #4
20718       \__fp_case_return_same_o:w
20719     \fi:
20720     \if_meaning:w #1 #4
20721       \__fp_case_return_o:Nw \c_zero_fp
20722     \fi:

```

```

20723 \if_meaning:w 2 #1
20724 \__fp_case_return_o:Nw \c_inf_fp
20725 \fi:
20726 \if_meaning:w 2 #3
20727 \__fp_case_return_o:Nw \c_inf_fp
20728 \else:
20729 \__fp_case_use:nw
20730 {
20731 \__fp_division_by_zero_o:NNww \c_inf_fp ^
20732 \s__fp \__fp_chk:w #1 #2 ;
20733 }
20734 \fi:
20735 \s__fp \__fp_chk:w #3#4
20736 }

```

(End definition for __fp_pow_zero_or_inf:ww.)

__fp_pow_normal_o:ww We have in front of us a , and $b \neq 0$, we know that a is a normal number, and we wish to compute $|a|^b$. If $|a| = 1$, we return 1, unless $a = -1$ and b is `nan`. Indeed, returning 1 at this point would wrongly raise “invalid” when the sign is considered. If $|a| \neq 1$, test the type of b :

- 0 Impossible, we already filtered $b = \pm 0$.
- 1 Call __fp_pow_npos_o:Nww.
- 2 Return $+\infty$ or $+0$ depending on the sign of b and whether the exponent of a is positive or not.
- 3 Return b .

```

20737 \cs_new:Npn \__fp_pow_normal_o:ww
20738 \s__fp \__fp_chk:w 1 #1#2#3; \s__fp \__fp_chk:w #4#5
20739 {
20740 \if_int_compare:w \__fp_str_if_eq:nn { #2 #3 }
20741 { 1 {1000} {0000} {0000} {0000} } = 0 \exp_stop_f:
20742 \if_int_compare:w #4 #1 = 32 \exp_stop_f:
20743 \exp_after:wN \__fp_case_return_ii_o:ww
20744 \fi:
20745 \__fp_case_return_o:Nww \c_one_fp
20746 \fi:
20747 \if_case:w #4 \exp_stop_f:
20748 \or:
20749 \exp_after:wN \__fp_pow_npos_o:Nww
20750 \exp_after:wN #5
20751 \or:
20752 \if_meaning:w 2 #5 \exp_after:wN \reverse_if:N \fi:
20753 \if_int_compare:w #2 > 0 \exp_stop_f:
20754 \exp_after:wN \__fp_case_return_o:Nww
20755 \exp_after:wN \c_inf_fp
20756 \else:
20757 \exp_after:wN \__fp_case_return_o:Nww
20758 \exp_after:wN \c_zero_fp
20759 \fi:
20760 \or:

```

```

20761     \__fp_case_return_ii_o:ww
20762 \fi:
20763 \s__fp \__fp_chk:w 1 #1 {#2} #3 ;
20764 \s__fp \__fp_chk:w #4 #5
20765 }

```

(End definition for __fp_pow_normal_o:ww.)

__fp_pow_npos_o:Nww We now know that $a \neq \pm 1$ is a normal number, and b is a normal number too. We want to compute $|a|^b = (|x| \cdot 10^n)^{y \cdot 10^p} = \exp((\ln|x| + n \ln(10)) \cdot y \cdot 10^p) = \exp(z)$. To compute the exponential accurately, we need to know the digits of z up to the 16-th position. Since the exponential of 10^5 is infinite, we only need at most 21 digits, hence the fixed point result of __fp_ln_o:w is precise enough for our needs. Start an integer expression for the decimal exponent of $e^{|z|}$. If z is negative, negate that decimal exponent, and prepare to take the inverse when converting from the fixed point to the floating point result.

```

20766 \cs_new:Npn \__fp_pow_npos_o:Nww #1 \s__fp \__fp_chk:w 1#2#3
20767 {
20768   \exp_after:wN \__fp_sanitize:Nw
20769   \exp_after:wN 0
20770   \int_value:w
20771   \if:w #1 \if_int_compare:w #3 > 0 \exp_stop_f: 0 \else: 2 \fi:
20772   \exp_after:wN \__fp_pow_npos_aux:NNnww
20773   \exp_after:wN +
20774   \exp_after:wN \__fp_fixed_to_float_o:wN
20775   \else:
20776   \exp_after:wN \__fp_pow_npos_aux:NNnww
20777   \exp_after:wN -
20778   \exp_after:wN \__fp_fixed_inv_to_float_o:wN
20779   \fi:
20780   {#3}
20781 }

```

(End definition for __fp_pow_npos_o:Nww.)

__fp_pow_npos_aux:NNnww The first argument is the conversion function from fixed point to float. Then comes an exponent and the 4 brace groups of x , followed by b . Compute $-\ln(x)$.

```

20782 \cs_new:Npn \__fp_pow_npos_aux:NNnww #1#2#3#4#5; \s__fp \__fp_chk:w 1#6#7#8;
20783 {
20784   #1
20785   \__fp_int_eval:w
20786   \__fp_ln_significand:NNNNnnnnN #4#5
20787   \__fp_pow_exponent:wnN {#3}
20788   \__fp_fixed_mul:wwn #8 {0000}{0000} ;
20789   \__fp_pow_B:wwN #7;
20790   #1 #2 0 % fixed_to_float_o:wN
20791 }
20792 \cs_new:Npn \__fp_pow_exponent:wnN #1; #2
20793 {
20794   \if_int_compare:w #2 > 0 \exp_stop_f:
20795   \exp_after:wN \__fp_pow_exponent:Nwnnnnnw % n\ln(10) - (-\ln(x))
20796   \exp_after:wN +
20797   \else:
20798   \exp_after:wN \__fp_pow_exponent:Nwnnnnnw % -(\ln|\ln(10) + (-\ln(x)))
20799   \exp_after:wN -

```



```

20800     \fi:
20801     #2; #1;
20802 }
20803 \cs_new:Npn \__fp_pow_exponent:Nwnnnnnw #1#2; #3#4#5#6#7#8;
20804 { %^A todo: use that in ln.
20805     \exp_after:wN \__fp_fixed_mul_after:wwn
20806     \int_value:w \__fp_int_eval:w \c__fp_leading_shift_int
20807     \exp_after:wN \__fp_pack:NNNNNw
20808     \int_value:w \__fp_int_eval:w \c__fp_middle_shift_int
20809     #1#2*23025 - #1 #3
20810     \exp_after:wN \__fp_pack:NNNNNw
20811     \int_value:w \__fp_int_eval:w \c__fp_middle_shift_int
20812     #1 #2*8509 - #1 #4
20813     \exp_after:wN \__fp_pack:NNNNNw
20814     \int_value:w \__fp_int_eval:w \c__fp_middle_shift_int
20815     #1 #2*2994 - #1 #5
20816     \exp_after:wN \__fp_pack:NNNNNw
20817     \int_value:w \__fp_int_eval:w \c__fp_middle_shift_int
20818     #1 #2*0456 - #1 #6
20819     \exp_after:wN \__fp_pack:NNNNNw
20820     \int_value:w \__fp_int_eval:w \c__fp_trailing_shift_int
20821     #1 #2*8401 - #1 #7
20822     #1 ( #2*7991 - #8 ) / 1 0000 ; ;
20823 }
20824 \cs_new:Npn \__fp_pow_B:wwN #1#2#3#4#5#6; #7;
20825 {
20826     \if_int_compare:w #7 < 0 \exp_stop_f:
20827     \exp_after:wN \__fp_pow_C_neg:w \int_value:w -
20828     \else:
20829     \if_int_compare:w #7 < 22 \exp_stop_f:
20830     \exp_after:wN \__fp_pow_C_pos:w \int_value:w
20831     \else:
20832     \exp_after:wN \__fp_pow_C_overflow:w \int_value:w
20833     \fi:
20834     \fi:
20835     #7 \exp_after:wN ;
20836     \int_value:w \__fp_int_eval:w 10 0000 + #1 \__fp_int_eval_end:
20837     #2#3#4#5#6 0000 0000 0000 0000 0000 0000 ; %^A todo: how many 0?
20838 }
20839 \cs_new:Npn \__fp_pow_C_overflow:w #1; #2; #3
20840 {
20841     + 2 * \c__fp_max_exponent_int
20842     \exp_after:wN \__fp_fixed_continue:wn \c__fp_one_fixed_tl
20843 }
20844 \cs_new:Npn \__fp_pow_C_neg:w #1 ; 1
20845 {
20846     \exp_after:wN \exp_after:wN \exp_after:wN \__fp_pow_C_pack:w
20847     \prg_replicate:nn {#1} {0}
20848 }
20849 \cs_new:Npn \__fp_pow_C_pos:w #1; 1
20850 { \__fp_pow_C_pos_loop:wN #1; }
20851 \cs_new:Npn \__fp_pow_C_pos_loop:wN #1; #2
20852 {
20853     \if_meaning:w 0 #1

```

```

20854     \exp_after:wN \__fp_pow_C_pack:w
20855     \exp_after:wN #2
20856   \else:
20857     \if_meaning:w 0 #2
20858     \exp_after:wN \__fp_pow_C_pos_loop:wN \int_value:w
20859   \else:
20860     \exp_after:wN \__fp_pow_C_overflow:w \int_value:w
20861   \fi:
20862   \__fp_int_eval:w #1 - 1 \exp_after:wN ;
20863 \fi:
20864 }
20865 \cs_new:Npn \__fp_pow_C_pack:w
20866 {
20867   \exp_after:wN \__fp_exp_large:NwN
20868   \exp_after:wN 5
20869   \c__fp_one_fixed_tl
20870 }

```

(End definition for __fp_pow_npos_aux:Nnnww.)

__fp_pow_neg:www
__fp_pow_neg_aux:wNN

This function is followed by three floating point numbers: a^b , $a \in [-\infty, -0]$, and b . If b is an even integer (case -1), $a^b = a^b$. If b is an odd integer (case 0), $a^b = -a^b$, obtained by a call to __fp_pow_neg_aux:wNN. Otherwise, the sign is undefined. This is invalid, unless a^b turns out to be +0 or nan, in which case we return that as a^b . In particular, since the underflow detection occurs before __fp_pow_neg:www is called, $(-0.1)**(12345.67)$ gives +0 rather than complaining that the sign is not defined.

```

20871 \cs_new:Npn \__fp_pow_neg:www \s__fp \__fp_chk:w #1#2; #3; #4;
20872 {
20873   \if_case:w \__fp_pow_neg_case:w #4 ;
20874     \exp_after:wN \__fp_pow_neg_aux:wNN
20875   \or:
20876     \if_int_compare:w \__fp_int_eval:w #1 / 2 = 1 \exp_stop_f:
20877     \__fp_invalid_operation_o:Nww ^ #3; #4;
20878     \exp:w \exp_end_continue_f:w
20879     \exp_after:wN \exp_after:wN
20880     \exp_after:wN \__fp_use_none_until_s:w
20881   \fi:
20882 \fi:
20883 \__fp_exp_after_o:w
20884 \s__fp \__fp_chk:w #1#2;
20885 }
20886 \cs_new:Npn \__fp_pow_neg_aux:wNN #1 \s__fp \__fp_chk:w #2#3
20887 {
20888   \exp_after:wN \__fp_exp_after_o:w
20889   \exp_after:wN \s__fp
20890   \exp_after:wN \__fp_chk:w
20891   \exp_after:wN #2
20892   \int_value:w \__fp_int_eval:w 2 - #3 \__fp_int_eval_end:
20893 }

```

(End definition for __fp_pow_neg:www and __fp_pow_neg_aux:wNN.)

__fp_pow_neg_case:w
__fp_pow_neg_case_aux:nnnnn
__fp_pow_neg_case_aux:Nnnw

This function expects a floating point number, and determines its “parity”. It should be used after \if_case:w or in an integer expression. It gives -1 if the number is an

even integer, 0 if the number is an odd integer, and 1 otherwise. Zeros and $\pm\infty$ are even (because very large finite floating points are even), while `nan` is a non-integer. The sign of normal numbers is irrelevant to parity. After `__fp_decimate:nNnnnn` the argument #1 of `__fp_pow_neg_case_aux:Nnnw` is a rounding digit, 0 if and only if the number was an integer, and #3 is the 8 least significant digits of that integer.

```

20894 \cs_new:Npn \__fp_pow_neg_case:w \s__fp \__fp_chk:w #1#2#3;
20895 {
20896   \if_case:w #1 \exp_stop_f:
20897     -1
20898   \or:   \__fp_pow_neg_case_aux:nnnnn #3
20899   \or:   -1
20900   \else: 1
20901   \fi:
20902   \exp_stop_f:
20903 }
20904 \cs_new:Npn \__fp_pow_neg_case_aux:nnnnn #1#2#3#4#5
20905 {
20906   \if_int_compare:w #1 > \c__fp_prec_int
20907     -1
20908   \else:
20909     \__fp_decimate:nNnnnn { \c__fp_prec_int - #1 }
20910     \__fp_pow_neg_case_aux:Nnnw
20911     {#2} {#3} {#4} {#5}
20912   \fi:
20913 }
20914 \cs_new:Npn \__fp_pow_neg_case_aux:Nnnw #1#2#3#4 ;
20915 {
20916   \if_meaning:w 0 #1
20917     \if_int_odd:w #3 \exp_stop_f:
20918     0
20919   \else:
20920     -1
20921   \fi:
20922   \else:
20923     1
20924   \fi:
20925 }

```

(End definition for `__fp_pow_neg_case:w`, `__fp_pow_neg_case_aux:nnnnn`, and `__fp_pow_neg_case_aux:Nnnw`.)

33.4 Factorial

`\c__fp_fact_max_arg_int` The maximum integer whose factorial fits in the exponent range is 3248, as $3249! \sim 10^{10000.8}$

```

20926 \int_const:Nn \c__fp_fact_max_arg_int { 3248 }

```

(End definition for `\c__fp_fact_max_arg_int`.)

`__fp_fact_o:w` First detect ± 0 and $+\infty$ and `nan`. Then note that factorial of anything with a negative sign (except -0) is undefined. Then call `__fp_small_int:wTF` to get an integer as the argument, and start a loop. This is not the most efficient way of computing the factorial,

but it works all right. Of course we work with 24 digits instead of 16. It is easy to check that computing factorials with this precision is enough.

```

20927 \cs_new:Npn \__fp_fact_o:w #1 \s_fp \__fp_chk:w #2#3#4; @
20928 {
20929   \if_case:w #2 \exp_stop_f:
20930     \__fp_case_return_o:Nw \c_one_fp
20931   \or:
20932   \or:
20933     \if_meaning:w 0 #3
20934     \exp_after:wN \__fp_case_return_same_o:w
20935   \fi:
20936   \or:
20937     \__fp_case_return_same_o:w
20938   \fi:
20939   \if_meaning:w 2 #3
20940     \__fp_case_use:nw { \__fp_invalid_operation_o:fw { fact } }
20941   \fi:
20942   \__fp_fact_pos_o:w
20943   \s_fp \__fp_chk:w #2 #3 #4 ;
20944 }

```

(End definition for __fp_fact_o:w.)

__fp_fact_pos_o:w Then check the input is an integer, and call __fp_facorial_int_o:n with that int as
 __fp_fact_int_o:w an argument. If it's too big the factorial overflows. Otherwise call __fp_sanitize:Nw
 with a positive sign marker 0 and an integer expression that will mop up any exponent
 in the calculation.

```

20945 \cs_new:Npn \__fp_fact_pos_o:w #1;
20946 {
20947   \__fp_small_int:wTF #1;
20948   { \__fp_fact_int_o:n }
20949   { \__fp_invalid_operation_o:fw { fact } #1; }
20950 }
20951 \cs_new:Npn \__fp_fact_int_o:n #1
20952 {
20953   \if_int_compare:w #1 > \c__fp_fact_max_arg_int
20954     \__fp_case_return:nw
20955     {
20956       \exp_after:wN \exp_after:wN \exp_after:wN \__fp_overflow:w
20957       \exp_after:wN \c_inf_fp
20958     }
20959   \fi:
20960   \exp_after:wN \__fp_sanitize:Nw
20961   \exp_after:wN 0
20962   \int_value:w \__fp_int_eval:w
20963   \__fp_fact_loop_o:w #1 . 4 , { 1 } { } { } { } { } { } { } ;
20964 }

```

(End definition for __fp_fact_pos_o:w and __fp_fact_int_o:w.)

__fp_fact_loop_o:w The loop receives an integer #1 whose factorial we want to compute, which we progres-
 sively decrement, and the result so far as an extended-precision number #2 in the form
 $\langle \text{exponent} \rangle, \langle \text{mantissa} \rangle$; The loop goes in steps of two because we compute $\#1 \cdot \#1 - 1$
 as an integer expression (it must fit since #1 is at most 3248), then multiply with the

result so far. We don't need to fill in most of the mantissa with zeros because `__fp_ep_mul:wwwn` first normalizes the extended precision number to avoid loss of precision. When reaching a small enough number simply use a table of factorials less than 10^8 . This limit is chosen because the normalization step cannot deal with larger integers.

```

20965 \cs_new:Npn \__fp_fact_loop_o:w #1 . #2 ;
20966 {
20967   \if_int_compare:w #1 < 12 \exp_stop_f:
20968     \__fp_fact_small_o:w #1
20969   \fi:
20970   \exp_after:wN \__fp_ep_mul:wwwn
20971   \exp_after:wN 4 \exp_after:wN ,
20972   \exp_after:wN { \int_value:w \__fp_int_eval:w #1 * (#1 - 1) }
20973   { } { } { } { } { } { } ;
20974   #2 ;
20975   {
20976     \exp_after:wN \__fp_fact_loop_o:w
20977     \int_value:w \__fp_int_eval:w #1 - 2 .
20978   }
20979 }
20980 \cs_new:Npn \__fp_fact_small_o:w #1 \fi: #2 ; #3 ; #4
20981 {
20982   \fi:
20983   \exp_after:wN \__fp_ep_mul:wwwn
20984   \exp_after:wN 4 \exp_after:wN ,
20985   \exp_after:wN
20986   {
20987     \int_value:w
20988     \if_case:w #1 \exp_stop_f:
20989       1 \or: 1 \or: 2 \or: 6 \or: 24 \or: 120 \or: 720 \or: 5040
20990       \or: 40320 \or: 362880 \or: 3628800 \or: 39916800
20991     \fi:
20992     } { } { } { } { } { } { } ;
20993     #3 ;
20994     \__fp_ep_to_float_o:wwN 0
20995   }
}

(End definition for \__fp_fact_loop_o:w.)

20996 \</package>

```

34 l3fp-trig Implementation

```

20997 \*package>
20998 \@@=fp>

\__fp_parse_word_acos:N Unary functions.
\__fp_parse_word_acosd:N
\__fp_parse_word_acsc:N
\__fp_parse_word_acscd:N
\__fp_parse_word_asec:N
\__fp_parse_word_asecd:N
\__fp_parse_word_asin:N
\__fp_parse_word_asind:N
\__fp_parse_word_cos:N
\__fp_parse_word_cosd:N
\__fp_parse_word_cot:N
\__fp_parse_word_cotd:N
\__fp_parse_word_csc:N
\__fp_parse_word_cscd:N
\__fp_parse_word_sec:N
\__fp_parse_word_secd:N
\__fp_parse_word_sin:N
\__fp_parse_word_sind:N

```

```

21006     {
21007         \exp_not:N \__fp_parse_unary_function:NNN
21008         \exp_not:c { __fp_#1_o:w }
21009         \exp_not:N \use_i:nn
21010     }
21011 \cs_new:cpx { __fp_parse_word_#1d:N }
21012 {
21013     \exp_not:N \__fp_parse_unary_function:NNN
21014     \exp_not:c { __fp_#1_o:w }
21015     \exp_not:N \use_ii:nn
21016 }
21017 }

```

(End definition for `__fp_parse_word_acos:N` and others.)

```

\__fp_parse_word_acot:N Those functions may receive a variable number of arguments.
\__fp_parse_word_acotd:N 21018 \cs_new:Npn \__fp_parse_word_acot:N
\__fp_parse_word_atan:N 21019 { \__fp_parse_function:NNN \__fp_acot_o:Nw \use_i:nn }
\__fp_parse_word_atand:N 21020 \cs_new:Npn \__fp_parse_word_acotd:N
21021 { \__fp_parse_function:NNN \__fp_acot_o:Nw \use_ii:nn }
21022 \cs_new:Npn \__fp_parse_word_atan:N
21023 { \__fp_parse_function:NNN \__fp_atan_o:Nw \use_i:nn }
21024 \cs_new:Npn \__fp_parse_word_atand:N
21025 { \__fp_parse_function:NNN \__fp_atan_o:Nw \use_ii:nn }

```

(End definition for `__fp_parse_word_acot:N` and others.)

34.1 Direct trigonometric functions

The approach for all trigonometric functions (sine, cosine, tangent, cotangent, cosecant, and secant), with arguments given in radians or in degrees, is the same.

- Filter out special cases (± 0 , $\pm \infty$ and NaN).
- Keep the sign for later, and work with the absolute value $|x|$ of the argument.
- Small numbers ($|x| < 1$ in radians, $|x| < 10$ in degrees) are converted to fixed point numbers (and to radians if $|x|$ is in degrees).
- For larger numbers, we need argument reduction. Subtract a multiple of $\pi/2$ (in degrees, 90) to bring the number to the range to $[0, \pi/2)$ (in degrees, $[0, 90)$).
- Reduce further to $[0, \pi/4]$ (in degrees, $[0, 45]$) using $\sin x = \cos(\pi/2 - x)$, and when working in degrees, convert to radians.
- Use the appropriate power series depending on the octant $\lfloor \frac{x}{\pi/4} \rfloor \bmod 8$ (in degrees, the same formula with $\pi/4 \rightarrow 45$), the sign, and the function to compute.

34.1.1 Filtering special cases

`__fp_sin_o:w` This function, and its analogs for `cos`, `csc`, `sec`, `tan`, and `cot` instead of `sin`, are followed either by `\use_i:nn` and a float in radians or by `\use_ii:nn` and a float in degrees. The sine of ± 0 or NaN is the same float. The sine of $\pm \infty$ raises an invalid operation exception with the appropriate function name. Otherwise, call the `trig` function to perform argument reduction and if necessary convert the reduced argument to radians.

Then, `__fp_sin_series_o:NNwww` is called to compute the Taylor series: this function receives a sign `#3`, an initial octant of 0, and the function `__fp_ep_to_float_o:wwN` which converts the result of the series to a floating point directly rather than taking its inverse, since $\sin(x) = \#3 \sin|x|$.

```

21026 \cs_new:Npn __fp_sin_o:w #1 \s__fp __fp_chk:w #2#3#4; @
21027 {
21028   \if_case:w #2 \exp_stop_f:
21029     __fp_case_return_same_o:w
21030   \or: __fp_case_use:nw
21031     {
21032       __fp_trig:NNNNwn #1 __fp_sin_series_o:NNwww
21033       __fp_ep_to_float_o:wwN #3 0
21034     }
21035   \or: __fp_case_use:nw
21036     { __fp_invalid_operation_o:fw { #1 { sin } { sind } } }
21037   \else: __fp_case_return_same_o:w
21038   \fi:
21039   \s__fp __fp_chk:w #2 #3 #4;
21040 }

```

(End definition for `__fp_sin_o:w`.)

`__fp_cos_o:w` The cosine of ± 0 is 1. The cosine of $\pm\infty$ raises an invalid operation exception. The cosine of NaN is itself. Otherwise, the `trig` function reduces the argument to at most half a right-angle and converts if necessary to radians. We then call the same series as for sine, but using a positive sign 0 regardless of the sign of x , and with an initial octant of 2, because $\cos(x) = +\sin(\pi/2 + |x|)$.

```

21041 \cs_new:Npn __fp_cos_o:w #1 \s__fp __fp_chk:w #2#3; @
21042 {
21043   \if_case:w #2 \exp_stop_f:
21044     __fp_case_return_o:Nw \c_one_fp
21045   \or: __fp_case_use:nw
21046     {
21047       __fp_trig:NNNNwn #1 __fp_sin_series_o:NNwww
21048       __fp_ep_to_float_o:wwN 0 2
21049     }
21050   \or: __fp_case_use:nw
21051     { __fp_invalid_operation_o:fw { #1 { cos } { cosd } } }
21052   \else: __fp_case_return_same_o:w
21053   \fi:
21054   \s__fp __fp_chk:w #2 #3;
21055 }

```

(End definition for `__fp_cos_o:w`.)

`__fp_csc_o:w` The cosecant of ± 0 is $\pm\infty$ with the same sign, with a division by zero exception (see `__fp_cot_zero_o:Nfw` defined below), which requires the function name. The cosecant of $\pm\infty$ raises an invalid operation exception. The cosecant of NaN is itself. Otherwise, the `trig` function performs the argument reduction, and converts if necessary to radians before calling the same series as for sine, using the sign `#3`, a starting octant of 0, and inverting during the conversion from the fixed point sine to the floating point result, because $\csc(x) = \#3(\sin|x|)^{-1}$.

```

21056 \cs_new:Npn __fp_csc_o:w #1 \s__fp __fp_chk:w #2#3#4; @

```

```

21057 {
21058     \if_case:w #2 \exp_stop_f:
21059         \__fp_cot_zero_o:Nfw #3 { #1 { csc } { cscd } }
21060     \or: \__fp_case_use:nw
21061         {
21062             \__fp_trig:NNNNNwn #1 \__fp_sin_series_o:NNwww
21063             \__fp_ep_inv_to_float_o:wwN #3 0
21064         }
21065     \or: \__fp_case_use:nw
21066         { \__fp_invalid_operation_o:fw { #1 { csc } { cscd } } }
21067     \else: \__fp_case_return_same_o:w
21068     \fi:
21069     \s__fp \__fp_chk:w #2 #3 #4;
21070 }

```

(End definition for __fp_csc_o:w.)

__fp_sec_o:w The secant of ± 0 is 1. The secant of $\pm\infty$ raises an invalid operation exception. The secant of NaN is itself. Otherwise, the `trig` function reduces the argument and turns it to radians before calling the same series as for sine, using a positive sign 0, a starting octant of 2, and inverting upon conversion, because $\sec(x) = +1/\sin(\pi/2 + |x|)$.

```

21071 \cs_new:Npn \__fp_sec_o:w #1 \s__fp \__fp_chk:w #2#3; @
21072 {
21073     \if_case:w #2 \exp_stop_f:
21074         \__fp_case_return_o:Nw \c_one_fp
21075     \or: \__fp_case_use:nw
21076         {
21077             \__fp_trig:NNNNNwn #1 \__fp_sin_series_o:NNwww
21078             \__fp_ep_inv_to_float_o:wwN 0 2
21079         }
21080     \or: \__fp_case_use:nw
21081         { \__fp_invalid_operation_o:fw { #1 { sec } { secd } } }
21082     \else: \__fp_case_return_same_o:w
21083     \fi:
21084     \s__fp \__fp_chk:w #2 #3;
21085 }

```

(End definition for __fp_sec_o:w.)

__fp_tan_o:w The tangent of ± 0 or NaN is the same floating point number. The tangent of $\pm\infty$ raises an invalid operation exception. Once more, the `trig` function does the argument reduction step and conversion to radians before calling `__fp_tan_series_o:NNwww`, with a sign #3 and an initial octant of 1 (this shift is somewhat arbitrary). See `__fp_cot_o:w` for an explanation of the 0 argument.

```

21086 \cs_new:Npn \__fp_tan_o:w #1 \s__fp \__fp_chk:w #2#3#4; @
21087 {
21088     \if_case:w #2 \exp_stop_f:
21089         \__fp_case_return_same_o:w
21090     \or: \__fp_case_use:nw
21091         {
21092             \__fp_trig:NNNNNwn #1
21093             \__fp_tan_series_o:NNwww 0 #3 1
21094         }
21095     \or: \__fp_case_use:nw

```



```

21096         { \__fp_invalid_operation_o:fw { #1 { tan } { tand } } }
21097     \else: \__fp_case_return_same_o:w
21098     \fi:
21099     \s__fp \__fp_chk:w #2 #3 #4;
21100 }

```

(End definition for __fp_tan_o:w.)

__fp_cot_o:w The cotangent of ± 0 is $\pm\infty$ with the same sign, with a division by zero exception (see
__fp_cot_zero_o:Nfw __fp_cot_zero_o:Nfw. The cotangent of $\pm\infty$ raises an invalid operation exception.
The cotangent of NaN is itself. We use $\cot x = -\tan(\pi/2 + x)$, and the initial octant
for the tangent was chosen to be 1, so the octant here starts at 3. The change in sign
is obtained by feeding __fp_tan_series_o:NNwww two signs rather than just the sign
of the argument: the first of those indicates whether we compute tangent or cotangent.
Those signs are eventually combined.

```

21101 \cs_new:Npn \__fp_cot_o:w #1 \s__fp \__fp_chk:w #2#3#4; @
21102 {
21103     \if_case:w #2 \exp_stop_f:
21104         \__fp_cot_zero_o:Nfw #3 { #1 { cot } { cotd } }
21105     \or: \__fp_case_use:nw
21106         {
21107             \__fp_trig:NNNNwn #1
21108             \__fp_tan_series_o:NNwww 2 #3 3
21109         }
21110     \or: \__fp_case_use:nw
21111         { \__fp_invalid_operation_o:fw { #1 { cot } { cotd } } }
21112     \else: \__fp_case_return_same_o:w
21113     \fi:
21114     \s__fp \__fp_chk:w #2 #3 #4;
21115 }
21116 \cs_new:Npn \__fp_cot_zero_o:Nfw #1#2#3 \fi:
21117 {
21118     \fi:
21119     \token_if_eq_meaning:NNTF 0 #1
21120     { \exp_args:NNf \__fp_division_by_zero_o:Nnw \c_inf_fp }
21121     { \exp_args:NNf \__fp_division_by_zero_o:Nnw \c_minus_inf_fp }
21122     {#2}
21123 }

```

(End definition for __fp_cot_o:w and __fp_cot_zero_o:Nfw.)

34.1.2 Distinguishing small and large arguments

__fp_trig:NNNNwn The first argument is \use_i:nn if the operand is in radians and \use_ii:nn if it is in
degrees. Arguments #2 to #5 control what trigonometric function we compute, and #6
to #8 are pieces of a normal floating point number. Call the _series function #2, with
arguments #3, either a conversion function (__fp_ep_to_float_o:wN or __fp_ep_-
inv_to_float_o:wN) or a sign 0 or 2 when computing tangent or cotangent; #4, a sign
0 or 2; the octant, computed in an integer expression starting with #5 and stopped by a
period; and a fixed point number obtained from the floating point number by argument
reduction (if necessary) and conversion to radians (if necessary). Any argument reduction
adjusts the octant accordingly by leaving a (positive) shift into its integer expression. Let
us explain the integer comparison. Two of the four \exp_after:wN are expanded, the

expansion hits the test, which is true if the float is at least 1 when working in radians, and at least 10 when working in degrees. Then one of the remaining `\exp_after:wN` hits `#1`, which picks the `trig` or `trigd` function in whichever branch of the conditional was taken. The final `\exp_after:wN` closes the conditional. At the end of the day, a number is `large` if it is ≥ 1 in radians or ≥ 10 in degrees, and `small` otherwise. All four `trig/trigd` auxiliaries receive the operand as an extended-precision number.

```

21124 \cs_new:Npn \__fp_trig:NNNNNwn #1#2#3#4#5 \s__fp \__fp_chk:w 1#6#7#8;
21125 {
21126   \exp_after:wN #2
21127   \exp_after:wN #3
21128   \exp_after:wN #4
21129   \int_value:w \__fp_int_eval:w #5
21130   \exp_after:wN \exp_after:wN \exp_after:wN \exp_after:wN
21131   \if_int_compare:w #7 > #1 0 1 \exp_stop_f:
21132   #1 \__fp_trig_large:ww \__fp_trigd_large:ww
21133   \else:
21134   #1 \__fp_trig_small:ww \__fp_trigd_small:ww
21135   \fi:
21136   #7,#8{0000}{0000};
21137 }

```

(End definition for `__fp_trig:NNNNNwn`.)

34.1.3 Small arguments

`__fp_trig_small:ww` This receives a small extended-precision number in radians and converts it to a fixed point number. Some trailing digits may be lost in the conversion, so we keep the original floating point number around: when computing sine or tangent (or their inverses), the last step is to multiply by the floating point number (as an extended-precision number) rather than the fixed point number. The period serves to end the integer expression for the octant.

```

21138 \cs_new:Npn \__fp_trig_small:ww #1,#2;
21139 { \__fp_ep_to_fixed:wwn #1,#2; . #1,#2; }

```

(End definition for `__fp_trig_small:ww`.)

`__fp_trigd_small:ww` Convert the extended-precision number to radians, then call `__fp_trig_small:ww` to massage it in the form appropriate for the `_series` auxiliary.

```

21140 \cs_new:Npn \__fp_trigd_small:ww #1,#2;
21141 {
21142   \__fp_ep_mul_raw:wwwN
21143   -1,{1745}{3292}{5199}{4329}{5769}{2369}; #1,#2;
21144   \__fp_trig_small:ww
21145 }

```

(End definition for `__fp_trigd_small:ww`.)

34.1.4 Argument reduction in degrees

`__fp_trigd_large:ww` Note that $25 \times 360 = 9000$, so $10^{k+1} \equiv 10^k \pmod{360}$ for $k \geq 3$. When the exponent `#1` is very large, we can thus safely replace it by 22 (or even 19). We turn the floating point number into a fixed point number with two blocks of 8 digits followed by five blocks of 4 digits. The original float is $100 \times \langle block_1 \rangle \cdots \langle block_3 \rangle . \langle block_4 \rangle \cdots \langle block_7 \rangle$, or is equal to

`__fp_trig_large_auxi:nnnnwNNNN`

`__fp_trigd_large_auxii:wNw`

`__fp_trigd_large_auxiii:www`

it modulo 360 if the exponent #1 is very large. The first auxiliary finds $\langle block_1 \rangle + \langle block_2 \rangle$ (mod 9), a single digit, and prepends it to the 4 digits of $\langle block_3 \rangle$. It also unpacks $\langle block_4 \rangle$ and grabs the 4 digits of $\langle block_7 \rangle$. The second auxiliary grabs the $\langle block_3 \rangle$ plus any contribution from the first two blocks as #1, the first digit of $\langle block_4 \rangle$ (just after the decimal point in hundreds of degrees) as #2, and the three other digits as #3. It finds the quotient and remainder of #1#2 modulo 9, adds twice the quotient to the integer expression for the octant, and places the remainder (between 0 and 8) before #3 to form a new $\langle block_4 \rangle$. The resulting fixed point number is $x \in [0, 0.9]$. If $x \geq 0.45$, we add 1 to the octant and feed $0.9 - x$ with an exponent of 2 (to compensate the fact that we are working in units of hundreds of degrees rather than degrees) to `_fp_trigd_small:ww`. Otherwise, we feed it x with an exponent of 2. The third auxiliary also discards digits which were not packed into the various $\langle blocks \rangle$. Since the original exponent #1 is at least 2, those are all 0 and no precision is lost (#6 and #7 are four 0 each).

```

21146 \cs_new:Npn \_fp_trigd_large:ww #1, #2#3#4#5#6#7;
21147 {
21148   \exp_after:wN \_fp_pack_eight:wNNNNNNNN
21149   \exp_after:wN \_fp_pack_eight:wNNNNNNNN
21150   \exp_after:wN \_fp_pack_twice_four:wNNNNNNNN
21151   \exp_after:wN \_fp_pack_twice_four:wNNNNNNNN
21152   \exp_after:wN \_fp_trigd_large_auxi:nnnnwNNNN
21153   \exp_after:wN ;
21154   \exp:w \exp_end_continue_f:w
21155   \prg_replicate:nn { \int_max:nn { 22 - #1 } { 0 } } { 0 }
21156   #2#3#4#5#6#7 0000 0000 0000 !
21157 }
21158 \cs_new:Npn \_fp_trigd_large_auxi:nnnnwNNNN #1#2#3#4#5; #6#7#8#9
21159 {
21160   \exp_after:wN \_fp_trigd_large_auxii:wNw
21161   \int_value:w \_fp_int_eval:w #1 + #2
21162   - (#1 + #2 - 4) / 9 * 9 \_fp_int_eval_end:
21163   #3;
21164   #4; #5{#6#7#8#9};
21165 }
21166 \cs_new:Npn \_fp_trigd_large_auxii:wNw #1; #2#3;
21167 {
21168   + (#1#2 - 4) / 9 * 2
21169   \exp_after:wN \_fp_trigd_large_auxiii:www
21170   \int_value:w \_fp_int_eval:w #1#2
21171   - (#1#2 - 4) / 9 * 9 \_fp_int_eval_end: #3 ;
21172 }
21173 \cs_new:Npn \_fp_trigd_large_auxiii:www #1; #2; #3!
21174 {
21175   \if_int_compare:w #1 < 4500 \exp_stop_f:
21176   \exp_after:wN \_fp_use_i_until_s:nw
21177   \exp_after:wN \_fp_fixed_continue:wn
21178   \else:
21179     + 1
21180   \fi:
21181   \_fp_fixed_sub:wnn {9000}{0000}{0000}{0000}{0000}{0000};
21182   {#1}#2{0000}{0000};
21183   { \_fp_trigd_small:ww 2, }
21184 }

```

(End definition for `_fp_trigd_large:ww` and others.)

34.1.5 Argument reduction in radians

Arguments greater or equal to 1 need to be reduced to a range where we only need a few terms of the Taylor series. We reduce to the range $[0, 2\pi]$ by subtracting multiples of 2π , then to the smaller range $[0, \pi/2]$ by subtracting multiples of $\pi/2$ (keeping track of how many times $\pi/2$ is subtracted), then to $[0, \pi/4]$ by mapping $x \rightarrow \pi/2 - x$ if appropriate. When the argument is very large, say, 10^{100} , an equally large multiple of 2π must be subtracted, hence we must work with a very good approximation of 2π in order to get a sensible remainder modulo 2π .

Specifically, we multiply the argument by an approximation of $1/(2\pi)$ with 10048 digits, then discard the integer part of the result, keeping 52 digits of the fractional part. From the fractional part of $x/(2\pi)$ we deduce the octant (quotient of the first three digits by 125). We then multiply by 8 or -8 (the latter when the octant is odd), ignore any integer part (related to the octant), and convert the fractional part to an extended precision number, before multiplying by $\pi/4$ to convert back to a value in radians in $[0, \pi/4]$.

It is possible to prove that given the precision of floating points and their range of exponents, the 52 digits may start at most with 24 zeros. The 5 last digits are affected by carries from computations which are not done, hence we are left with at least $52 - 24 - 5 = 23$ significant digits, enough to round correctly up to $0.6 \cdot \text{ulp}$ in all cases.

`\c_fp_trig_intarray` This integer array stores blocks of 8 decimals of $10^{-16}/(2\pi)$. Each entry is 10^8 plus an 8 digit number storing 8 decimals. In total we store 10112 decimals of $10^{-16}/(2\pi)$. The number of decimals we really need is the maximum exponent plus the number of digits we later need, 52, plus 12 ($4 - 1$ groups of 4 digits). The memory footprint (1/2 byte per digit) is the same as an earlier method of storing the data as a control sequence name, but the major advantage is that we can unpack specific subsets of the digits without unpacking the 10112 decimals.

```

21185 \intarray_const_from_clist:Nn \c\_fp_trig_intarray
21186 {
21187     100000000, 100000000, 115915494, 130918953, 135768883, 176337251,
21188     143620344, 159645740, 145644874, 176673440, 158896797, 163422653,
21189     150901138, 102766253, 108595607, 128427267, 157958036, 189291184,
21190     161145786, 152877967, 141073169, 198392292, 139966937, 140907757,
21191     130777463, 196925307, 168871739, 128962173, 197661693, 136239024,
21192     117236290, 111832380, 111422269, 197557159, 140461890, 108690267,
21193     139561204, 189410936, 193784408, 155287230, 199946443, 140024867,
21194     123477394, 159610898, 132309678, 130749061, 166986462, 180469944,
21195     186521878, 181574786, 156696424, 110389958, 174139348, 160998386,
21196     180991999, 162442875, 158517117, 188584311, 117518767, 116054654,
21197     175369880, 109739460, 136475933, 137680593, 102494496, 163530532,
21198     171567755, 103220324, 177781639, 171660229, 146748119, 159816584,
21199     106060168, 103035998, 113391198, 174988327, 186654435, 127975507,
21200     100162406, 177564388, 184957131, 108801221, 199376147, 168137776,
21201     147378906, 133068046, 145797848, 117613124, 127314069, 196077502,
21202     145002977, 159857089, 105690279, 167851315, 125210016, 131774602,
21203     109248116, 106240561, 145620314, 164840892, 148459191, 143521157,
21204     154075562, 100871526, 160680221, 171591407, 157474582, 172259774,
21205     162853998, 175155329, 139081398, 117724093, 158254797, 107332871,
21206     190406999, 175907657, 170784934, 170393589, 182808717, 134256403,

```

21207	166895116,	162545705,	194332763,	112686500,	126122717,	197115321,
21208	112599504,	138667945,	103762556,	108363171,	116952597,	158128224,
21209	194162333,	143145106,	112353687,	185631136,	136692167,	114206974,
21210	169601292,	150578336,	105311960,	185945098,	139556718,	170995474,
21211	165104316,	123815517,	158083944,	129799709,	199505254,	138756612,
21212	194458833,	106846050,	178529151,	151410404,	189298850,	163881607,
21213	176196993,	107341038,	199957869,	118905980,	193737772,	106187543,
21214	122271893,	101366255,	126123878,	103875388,	181106814,	106765434,
21215	108282785,	126933426,	179955607,	107903860,	160352738,	199624512,
21216	159957492,	176297023,	159409558,	143011648,	129641185,	157771240,
21217	157544494,	157021789,	176979240,	194903272,	194770216,	164960356,
21218	153181535,	144003840,	168987471,	176915887,	163190966,	150696440,
21219	147769706,	187683656,	177810477,	197954503,	153395758,	130188183,
21220	186879377,	166124814,	195305996,	155802190,	183598751,	103512712,
21221	190432315,	180498719,	168687775,	194656634,	162210342,	104440855,
21222	149785037,	192738694,	129353661,	193778292,	187359378,	143470323,
21223	102371458,	137923557,	111863634,	119294601,	183182291,	196416500,
21224	187830793,	131353497,	179099745,	186492902,	167450609,	189368909,
21225	145883050,	133703053,	180547312,	132158094,	131976760,	132283131,
21226	141898097,	149822438,	133517435,	169898475,	101039500,	168388003,
21227	197867235,	199608024,	100273901,	108749548,	154787923,	156826113,
21228	199489032,	168997427,	108349611,	149208289,	103776784,	174303550,
21229	145684560,	183671479,	130845672,	133270354,	185392556,	120208683,
21230	193240995,	162211753,	131839402,	109707935,	170774965,	149880868,
21231	160663609,	168661967,	103747454,	121028312,	119251846,	122483499,
21232	111611495,	166556037,	196967613,	199312829,	196077608,	127799010,
21233	107830360,	102338272,	198790854,	102387615,	157445430,	192601191,
21234	100543379,	198389046,	154921248,	129516070,	172853005,	122721023,
21235	160175233,	113173179,	175931105,	103281551,	109373913,	163964530,
21236	157926071,	180083617,	195487672,	146459804,	173977292,	144810920,
21237	109371257,	186918332,	189588628,	139904358,	168666639,	175673445,
21238	114095036,	137327191,	174311388,	106638307,	125923027,	159734506,
21239	105482127,	178037065,	133778303,	121709877,	134966568,	149080032,
21240	169885067,	141791464,	168350828,	116168533,	114336160,	173099514,
21241	198531198,	119733758,	144420984,	116559541,	152250643,	139431286,
21242	144403838,	183561508,	179771645,	101706470,	167518774,	156059160,
21243	187168578,	157939226,	123475633,	117111329,	198655941,	159689071,
21244	198506887,	144230057,	151919770,	156900382,	118392562,	120338742,
21245	135362568,	108354156,	151729710,	188117217,	195936832,	156488518,
21246	174997487,	108553116,	159830610,	113921445,	144601614,	188452770,
21247	125114110,	170248521,	173974510,	138667364,	103872860,	109967489,
21248	131735618,	112071174,	104788993,	168886556,	192307848,	150230570,
21249	157144063,	163863202,	136852010,	174100574,	185922811,	115721968,
21250	100397824,	175953001,	166958522,	112303464,	118773650,	143546764,
21251	164565659,	171901123,	108476709,	193097085,	191283646,	166919177,
21252	169387914,	133315566,	150669813,	121641521,	100895711,	172862384,
21253	126070678,	145176011,	113450800,	169947684,	122356989,	162488051,
21254	157759809,	153397080,	185475059,	175362656,	149034394,	145420581,
21255	178864356,	183042000,	131509559,	147434392,	152544850,	167491429,
21256	108647514,	142303321,	133245695,	111634945,	167753939,	142403609,
21257	105438335,	152829243,	142203494,	184366151,	146632286,	102477666,
21258	166049531,	140657343,	157553014,	109082798,	180914786,	169343492,
21259	127376026,	134997829,	195701816,	119643212,	133140475,	176289748,
21260	140828911,	174097478,	126378991,	181699939,	148749771,	151989818,

21261	172666294,	160183053,	195832752,	109236350,	168538892,	128468247,
21262	125997252,	183007668,	156937583,	165972291,	198244297,	147406163,
21263	181831139,	158306744,	134851692,	185973832,	137392662,	140243450,
21264	119978099,	140402189,	161348342,	173613676,	144991382,	171541660,
21265	163424829,	136374185,	106122610,	186132119,	198633462,	184709941,
21266	183994274,	129559156,	128333990,	148038211,	175011612,	111667205,
21267	119125793,	103552929,	124113440,	131161341,	112495318,	138592695,
21268	184904438,	146807849,	109739828,	108855297,	104515305,	139914009,
21269	188698840,	188365483,	166522246,	168624087,	125401404,	100911787,
21270	142122045,	123075334,	173972538,	114940388,	141905868,	142311594,
21271	163227443,	139066125,	116239310,	162831953,	123883392,	113153455,
21272	163815117,	152035108,	174595582,	101123754,	135976815,	153401874,
21273	107394340,	136339780,	138817210,	104531691,	182951948,	179591767,
21274	139541778,	179243527,	161740724,	160593916,	102732282,	187946819,
21275	136491289,	149714953,	143255272,	135916592,	198072479,	198580612,
21276	169007332,	118844526,	179433504,	155801952,	149256630,	162048766,
21277	116134365,	133992028,	175452085,	155344144,	109905129,	182727454,
21278	165911813,	122232840,	151166615,	165070983,	175574337,	129548631,
21279	120411217,	116380915,	160616116,	157320000,	183306114,	160618128,
21280	103262586,	195951602,	146321661,	138576614,	180471993,	127077713,
21281	116441201,	159496011,	106328305,	120759583,	148503050,	179095584,
21282	198298218,	167402898,	138551383,	123957020,	180763975,	150429225,
21283	198476470,	171016426,	197438450,	143091658,	164528360,	132493360,
21284	143546572,	137557916,	113663241,	120457809,	196971566,	134022158,
21285	180545794,	131328278,	100552461,	132088901,	187421210,	192448910,
21286	141005215,	149680971,	113720754,	100571096,	134066431,	135745439,
21287	191597694,	135788920,	179342561,	177830222,	137011486,	142492523,
21288	192487287,	113132021,	176673607,	156645598,	127260957,	141566023,
21289	143787436,	129132109,	174858971,	150713073,	191040726,	143541417,
21290	197057222,	165479803,	181512759,	157912400,	125344680,	148220261,
21291	173422990,	101020483,	106246303,	137964746,	178190501,	181183037,
21292	151538028,	179523433,	141955021,	135689770,	191290561,	143178787,
21293	192086205,	174499925,	178975690,	118492103,	124206471,	138519113,
21294	188147564,	102097605,	154895793,	178514140,	141453051,	151583964,
21295	128232654,	106020603,	131189158,	165702720,	186250269,	191639375,
21296	115278873,	160608114,	155694842,	110322407,	177272742,	116513642,
21297	134366992,	171634030,	194053074,	180652685,	109301658,	192136921,
21298	141431293,	171341061,	157153714,	106203978,	147618426,	150297807,
21299	186062669,	169960809,	118422347,	163350477,	146719017,	145045144,
21300	161663828,	146208240,	186735951,	102371302,	190444377,	194085350,
21301	134454426,	133413062,	163074595,	113830310,	122931469,	134466832,
21302	185176632,	182415152,	110179422,	164439571,	181217170,	121756492,
21303	119644493,	196532222,	118765848,	182445119,	109401340,	150443213,
21304	198586286,	121083179,	139396084,	143898019,	114787389,	177233102,
21305	186310131,	148695521,	126205182,	178063494,	157118662,	177825659,
21306	188310053,	151552316,	165984394,	109022180,	163144545,	121212978,
21307	197344714,	188741258,	126822386,	102360271,	109981191,	152056882,
21308	134723983,	158013366,	106837863,	128867928,	161973236,	172536066,
21309	185216856,	132011948,	197807339,	158419190,	166595838,	167852941,
21310	124187182,	117279875,	106103946,	106481958,	157456200,	160892122,
21311	184163943,	173846549,	158993202,	184812364,	133466119,	170732430,
21312	195458590,	173361878,	162906318,	150165106,	126757685,	112163575,
21313	188696307,	145199922,	100107766,	176830946,	198149756,	122682434,
21314	179367131,	108412102,	119520899,	148191244,	140487511,	171059184,

21315	141399078,	189455775,	118462161,	190415309,	134543802,	180893862,
21316	180732375,	178615267,	179711433,	123241969,	185780563,	176301808,
21317	184386640,	160717536,	183213626,	129671224,	126094285,	140110963,
21318	121826276,	151201170,	122552929,	128965559,	146082049,	138409069,
21319	107606920,	103954646,	119164002,	115673360,	117909631,	187289199,
21320	186343410,	186903200,	157966371,	103128612,	135698881,	176403642,
21321	152540837,	109810814,	183519031,	121318624,	172281810,	150845123,
21322	169019064,	166322359,	138872454,	163073727,	128087898,	130041018,
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21324	196315803,	186381054,	124578934,	150084553,	128031351,	118843410,
21325	107373060,	159565443,	173624887,	171292628,	198074235,	139074061,
21326	178690578,	144431052,	174262641,	176783005,	182214864,	162289361,
21327	192966929,	192033046,	169332843,	181580535,	164864073,	118444059,
21328	195496893,	153773183,	167266131,	130108623,	158802128,	180432893,
21329	144562140,	147978945,	142337360,	158506327,	104399819,	132635916,
21330	168734194,	136567839,	101281912,	120281622,	195003330,	112236091,
21331	185875592,	101959081,	122415367,	194990954,	148881099,	175891989,
21332	108115811,	163538891,	163394029,	123722049,	184837522,	142362091,
21333	100834097,	156679171,	100841679,	157022331,	178971071,	102928884,
21334	189701309,	195339954,	124415335,	106062584,	139214524,	133864640,
21335	134324406,	157317477,	155340540,	144810061,	177612569,	108474646,
21336	114329765,	143900008,	138265211,	145210162,	136643111,	197987319,
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21338	102901215,	167964187,	135746835,	187317233,	110047459,	163339773,
21339	124770449,	118885134,	141536376,	100915375,	164267438,	145016622,
21340	113937193,	106748706,	128815954,	164819775,	119220771,	102367432,
21341	189062690,	170911791,	194127762,	112245117,	123546771,	115640433,
21342	135772061,	166615646,	174474627,	130562291,	133320309,	153340551,
21343	138417181,	194605321,	150142632,	180008795,	151813296,	175497284,
21344	167018836,	157425342,	150169942,	131069156,	134310662,	160434122,
21345	105213831,	158797111,	150754540,	163290657,	102484886,	148697402,
21346	187203725,	198692811,	149360627,	140384233,	128749423,	132178578,
21347	177507355,	171857043,	178737969,	134023369,	102911446,	196144864,
21348	197697194,	134527467,	144296030,	189437192,	154052665,	188907106,
21349	162062575,	150993037,	199766583,	167936112,	181374511,	104971506,
21350	115378374,	135795558,	167972129,	135876446,	130937572,	103221320,
21351	124605656,	161129971,	131027586,	191128460,	143251843,	143269155,
21352	129284585,	173495971,	150425653,	199302112,	118494723,	121323805,
21353	116549802,	190991967,	168151180,	122483192,	151273721,	199792134,
21354	133106764,	121874844,	126215985,	112167639,	167793529,	182985195,
21355	185453921,	106957880,	158685312,	132775454,	133229161,	198905318,
21356	190537253,	191582222,	192325972,	178133427,	181825606,	148823337,
21357	160719681,	101448145,	131983362,	137910767,	112550175,	128826351,
21358	183649210,	135725874,	110356573,	189469487,	154446940,	118175923,
21359	106093708,	128146501,	185742532,	149692127,	164624247,	183221076,
21360	154737505,	168198834,	156410354,	158027261,	125228550,	131543250,
21361	139591848,	191898263,	104987591,	115406321,	103542638,	190012837,
21362	142615518,	178773183,	175862355,	117537850,	169565995,	170028011,
21363	158412588,	170150030,	117025916,	174630208,	142412449,	112839238,
21364	105257725,	114737141,	123102301,	172563968,	130555358,	132628403,
21365	183638157,	168682846,	143304568,	105994018,	170010719,	152092970,
21366	117799058,	132164175,	179868116,	158654714,	177489647,	116547948,
21367	183121404,	131836079,	184431405,	157311793,	149677763,	173989893,
21368	102277656,	107058530,	140837477,	152640947,	143507039,	152145247,

```

21369      101683884, 107090870, 161471944, 137225650, 128231458, 172995869,
21370      173831689, 171268519, 139042297, 111072135, 107569780, 137262545,
21371      181410950, 138270388, 198736451, 162848201, 180468288, 120582913,
21372      153390138, 135649144, 130040157, 106509887, 192671541, 174507066,
21373      186888783, 143805558, 135011967, 145862340, 180595327, 124727843,
21374      182925939, 157715840, 136885940, 198993925, 152416883, 178793572,
21375      179679516, 154076673, 192703125, 164187609, 162190243, 104699348,
21376      159891990, 160012977, 174692145, 132970421, 167781726, 115178506,
21377      153008552, 155999794, 102099694, 155431545, 127458567, 104403686,
21378      168042864, 184045128, 181182309, 179349696, 127218364, 192935516,
21379      120298724, 169583299, 148193297, 183358034, 159023227, 105261254,
21380      121144370, 184359584, 194433836, 138388317, 175184116, 108817112,
21381      151279233, 137457721, 193398208, 119005406, 132929377, 175306906,
21382      160741530, 149976826, 147124407, 176881724, 186734216, 185881509,
21383      191334220, 175930947, 117385515, 193408089, 157124410, 163472089,
21384      131949128, 180783576, 131158294, 100549708, 191802336, 165960770,
21385      170927599, 101052702, 181508688, 197828549, 143403726, 142729262,
21386      110348701, 139928688, 153550062, 106151434, 130786653, 196085995,
21387      100587149, 139141652, 106530207, 100852656, 124074703, 166073660,
21388      153338052, 163766757, 120188394, 197277047, 122215363, 138511354,
21389      183463624, 161985542, 159938719, 133367482, 104220974, 149956672,
21390      170250544, 164232439, 157506869, 159133019, 137469191, 142980999,
21391      134242305, 150172665, 121209241, 145596259, 160554427, 159095199,
21392      168243130, 184279693, 171132070, 121049823, 123819574, 171759855,
21393      119501864, 163094029, 175943631, 194450091, 191506160, 149228764,
21394      132319212, 197034460, 193584259, 126727638, 168143633, 109856853,
21395      127860243, 132141052, 133076065, 188414958, 158718197, 107124299,
21396      159592267, 181172796, 144388537, 196763139, 127431422, 179531145,
21397      100064922, 112650013, 132686230, 121550837,
21398      }

```

(End definition for `\c__fp_trig_intarray`.)

`__fp_trig_large:ww` The exponent #1 is between 1 and 10000. We wish to look up decimals $10^{\#1-16}/(2\pi)$ starting from the digit #1 + 1. Since they are stored in batches of 8, compute $\lfloor \#1/8 \rfloor$ and fetch blocks of 8 digits starting there. The numbering of items in `\c__fp_trig_intarray` starts at 1, so the block $\lfloor \#1/8 \rfloor + 1$ contains the digit we want, at one of the eight positions. Each call to `\int_value:w __kernel_intarray_item:Nn` expands the next, until being stopped by `__fp_trig_large_auxiii:w` using `\exp_stop_f:.` Once all these blocks are unpacked, the `\exp_stop_f:` and 0 to 7 digits are removed by `\use_none:n...n`. Finally, `__fp_trig_large_auxii:w` packs 64 digits (there are between 65 and 72 at this point) into groups of 4 and the auxv auxiliary is called.

```

21399 \cs_new:Npn \__fp_trig_large:ww #1, #2#3#4#5#6;
21400 {
21401   \exp_after:wN \__fp_trig_large_auxi:w
21402   \int_value:w \__fp_int_eval:w (#1 - 4) / 8 \exp_after:wN ,
21403   \int_value:w #1 , ;
21404   {#2}{#3}{#4}{#5} ;
21405 }
21406 \cs_new:Npn \__fp_trig_large_auxi:w #1, #2,
21407 {
21408   \exp_after:wN \exp_after:wN
21409   \exp_after:wN \__fp_trig_large_auxii:w
21410   \cs:w

```



```

21411     use_none:n \prg_replicate:nn { #2 - #1 * 8 } { n }
21412     \exp_after:wN
21413   \cs_end:
21414   \int_value:w
21415   \__kernel_intarray_item:Nn \c__fp_trig_intarray
21416     { \__fp_int_eval:w #1 + 1 \scan_stop: }
21417   \exp_after:wN \__fp_trig_large_auxiii:w \int_value:w
21418   \__kernel_intarray_item:Nn \c__fp_trig_intarray
21419     { \__fp_int_eval:w #1 + 2 \scan_stop: }
21420   \exp_after:wN \__fp_trig_large_auxiii:w \int_value:w
21421   \__kernel_intarray_item:Nn \c__fp_trig_intarray
21422     { \__fp_int_eval:w #1 + 3 \scan_stop: }
21423   \exp_after:wN \__fp_trig_large_auxiii:w \int_value:w
21424   \__kernel_intarray_item:Nn \c__fp_trig_intarray
21425     { \__fp_int_eval:w #1 + 4 \scan_stop: }
21426   \exp_after:wN \__fp_trig_large_auxiii:w \int_value:w
21427   \__kernel_intarray_item:Nn \c__fp_trig_intarray
21428     { \__fp_int_eval:w #1 + 5 \scan_stop: }
21429   \exp_after:wN \__fp_trig_large_auxiii:w \int_value:w
21430   \__kernel_intarray_item:Nn \c__fp_trig_intarray
21431     { \__fp_int_eval:w #1 + 6 \scan_stop: }
21432   \exp_after:wN \__fp_trig_large_auxiii:w \int_value:w
21433   \__kernel_intarray_item:Nn \c__fp_trig_intarray
21434     { \__fp_int_eval:w #1 + 7 \scan_stop: }
21435   \exp_after:wN \__fp_trig_large_auxiii:w \int_value:w
21436   \__kernel_intarray_item:Nn \c__fp_trig_intarray
21437     { \__fp_int_eval:w #1 + 8 \scan_stop: }
21438   \exp_after:wN \__fp_trig_large_auxiii:w \int_value:w
21439   \__kernel_intarray_item:Nn \c__fp_trig_intarray
21440     { \__fp_int_eval:w #1 + 9 \scan_stop: }
21441   \exp_stop_f:
21442 }
21443 \cs_new:Npn \__fp_trig_large_auxii:w
21444 {
21445   \__fp_pack_twice_four:wNNNNNNNN \__fp_pack_twice_four:wNNNNNNNN
21446   \__fp_pack_twice_four:wNNNNNNNN \__fp_pack_twice_four:wNNNNNNNN
21447   \__fp_pack_twice_four:wNNNNNNNN \__fp_pack_twice_four:wNNNNNNNN
21448   \__fp_pack_twice_four:wNNNNNNNN \__fp_pack_twice_four:wNNNNNNNN
21449   \__fp_trig_large_auxv:www ;
21450 }
21451 \cs_new:Npn \__fp_trig_large_auxiii:w 1 { \exp_stop_f: }

```

(End definition for __fp_trig_large:ww and others.)

```

\__fp_trig_large_auxv:www
\__fp_trig_large_auxvi:wNNNNNNNN
\__fp_trig_large_pack:NNNNw

```

First come the first 64 digits of the fractional part of $10^{1-16}/(2\pi)$, arranged in 16 blocks of 4, and ending with a semicolon. Then a few more digits of the same fractional part, ending with a semicolon, then 4 blocks of 4 digits holding the significand of the original argument. Multiply the 16-digit significand with the 64-digit fractional part: the `auxvi` auxiliary receives the significand as `#2#3#4#5` and 16 digits of the fractional part as `#6#7#8#9`, and computes one step of the usual ladder of pack functions we use for multiplication (see *e.g.*, `__fp_fixed_mul:wN`), then discards one block of the fractional part to set things up for the next step of the ladder. We perform 13 such steps, replacing the last `middle` shift by the appropriate `trailing` shift, then discard the significand and remaining 3 blocks from the fractional part, as there are not enough digits to compute

any more step in the ladder. The last semicolon closes the ladder, and we return control to the `auxvii` auxiliary.

```

21452 \cs_new:Npn \__fp_trig_large_auxv:www #1; #2; #3;
21453 {
21454   \exp_after:wN \__fp_use_i_until_s:nw
21455   \exp_after:wN \__fp_trig_large_auxvii:w
21456   \int_value:w \__fp_int_eval:w \c__fp_leading_shift_int
21457   \prg_replicate:nn { 13 }
21458   { \__fp_trig_large_auxvi:wnnnnnnnn }
21459   + \c__fp_trailing_shift_int - \c__fp_middle_shift_int
21460   \__fp_use_i_until_s:nw
21461   ; #3 #1 ; ;
21462 }
21463 \cs_new:Npn \__fp_trig_large_auxvi:wnnnnnnnn #1; #2#3#4#5#6#7#8#9
21464 {
21465   \exp_after:wN \__fp_trig_large_pack:NNNNNw
21466   \int_value:w \__fp_int_eval:w \c__fp_middle_shift_int
21467   + #2*#9 + #3*#8 + #4*#7 + #5*#6
21468   #1; {#2}{#3}{#4}{#5} {#7}{#8}{#9}
21469 }
21470 \cs_new:Npn \__fp_trig_large_pack:NNNNNw #1#2#3#4#5#6;
21471 { + #1#2#3#4#5 ; #6 }

```

(End definition for `__fp_trig_large_auxv:www`, `__fp_trig_large_auxvi:wnnnnnnnn`, and `__fp_trig_large_pack:NNNNNw`.)

```

\__fp_trig_large_auxvii:w
\__fp_trig_large_auxviii:w
\__fp_trig_large_auxix:Nw
\__fp_trig_large_auxx:wNNNNN
\__fp_trig_large_auxxi:w

```

The `auxvii` auxiliary is followed by 52 digits and a semicolon. We find the octant as the integer part of 8 times what follows, or equivalently as the integer part of $\#1\#2\#3/125$, and add it to the surrounding integer expression for the octant. We then compute 8 times the 52-digit number, with a minus sign if the octant is odd. Again, the last `middle` shift is converted to a `trailing` shift. Any integer part (including negative values which come up when the octant is odd) is discarded by `__fp_use_i_until_s:nw`. The resulting fractional part should then be converted to radians by multiplying by $2\pi/8$, but first, build an extended precision number by abusing `__fp_ep_to_ep_loop:N` with the appropriate trailing markers. Finally, `__fp_trig_small:ww` sets up the argument for the functions which compute the Taylor series.

```

21472 \cs_new:Npn \__fp_trig_large_auxvii:w #1#2#3
21473 {
21474   \exp_after:wN \__fp_trig_large_auxviii:ww
21475   \int_value:w \__fp_int_eval:w (#1#2#3 - 62) / 125 ;
21476   #1#2#3
21477 }
21478 \cs_new:Npn \__fp_trig_large_auxviii:ww #1;
21479 {
21480   + #1
21481   \if_int_odd:w #1 \exp_stop_f:
21482     \exp_after:wN \__fp_trig_large_auxix:Nw
21483     \exp_after:wN -
21484   \else:
21485     \exp_after:wN \__fp_trig_large_auxix:Nw
21486     \exp_after:wN +
21487   \fi:
21488 }

```

```

21489 \cs_new:Npn \__fp_trig_large_auxix:Nw
21490 {
21491   \exp_after:wN \__fp_use_i_until_s:nw
21492   \exp_after:wN \__fp_trig_large_auxxi:w
21493   \int_value:w \__fp_int_eval:w \c__fp_leading_shift_int
21494   \prg_replicate:nn { 13 }
21495   { \__fp_trig_large_auxx:wNNNNN }
21496   + \c__fp_trailing_shift_int - \c__fp_middle_shift_int
21497   ;
21498 }
21499 \cs_new:Npn \__fp_trig_large_auxx:wNNNNN #1; #2 #3#4#5#6
21500 {
21501   \exp_after:wN \__fp_trig_large_pack:NNNNNw
21502   \int_value:w \__fp_int_eval:w \c__fp_middle_shift_int
21503   #2 8 * #3#4#5#6
21504   #1; #2
21505 }
21506 \cs_new:Npn \__fp_trig_large_auxxi:w #1;
21507 {
21508   \exp_after:wN \__fp_ep_mul_raw:wwwN
21509   \int_value:w \__fp_int_eval:w 0 \__fp_ep_to_ep_loop:N #1 ; ; !
21510   0,{7853}{9816}{3397}{4483}{0961}{5661};
21511   \__fp_trig_small:ww
21512 }

```

(End definition for __fp_trig_large_auxvii:w and others.)

34.1.6 Computing the power series

__fp_sin_series_o:NNwww Here we receive a conversion function __fp_ep_to_float_o:wwN or __fp_ep_inv_to_float_o:wwN, a *sign* (0 or 2), a (non-negative) *octant* delimited by a dot, a *fixed point* number delimited by a semicolon, and an extended-precision number. The auxiliary receives:

- the conversion function #1;
- the final sign, which depends on the octant #3 and the sign #2;
- the octant #3, which controls the series we use;
- the square #4 * #4 of the argument as a fixed point number, computed with __fp_fixed_mul:wwn;
- the number itself as an extended-precision number.

If the octant is in $\{1, 2, 5, 6, \dots\}$, we are near an extremum of the function and we use the series

$$\cos(x) = 1 - x^2 \left(\frac{1}{2!} - x^2 \left(\frac{1}{4!} - x^2 \left(\dots \right) \right) \right).$$

Otherwise, the series

$$\sin(x) = x \left(1 - x^2 \left(\frac{1}{3!} - x^2 \left(\frac{1}{5!} - x^2 \left(\dots \right) \right) \right) \right)$$

is used. Finally, the extended-precision number is converted to a floating point number with the given sign, and `__fp_sanitizew` checks for overflow and underflow.

```

21513 \cs_new:Npn \__fp_sin_series_o:NNwww #1#2#3. #4;
21514 {
21515   \__fp_fixed_mul:wwn #4; #4;
21516   {
21517     \exp_after:wN \__fp_sin_series_aux_o:NNwww
21518     \exp_after:wN #1
21519     \int_value:w
21520     \if_int_odd:w \__fp_int_eval:w (#3 + 2) / 4 \__fp_int_eval_end:
21521       #2
21522     \else:
21523       \if_meaning:w #2 0 2 \else: 0 \fi:
21524     \fi:
21525     {#3}
21526   }
21527 }
21528 \cs_new:Npn \__fp_sin_series_aux_o:NNwww #1#2#3 #4; #5,#6;
21529 {
21530   \if_int_odd:w \__fp_int_eval:w #3 / 2 \__fp_int_eval_end:
21531     \exp_after:wN \use_i:nn
21532   \else:
21533     \exp_after:wN \use_ii:nn
21534   \fi:
21535   { % 1/18!
21536     \__fp_fixed_mul_sub_back:wwwn {0000}{0000}{0000}{0001}{5619}{2070};
21537     #4;{0000}{0000}{0000}{0477}{9477}{3324};
21538     \__fp_fixed_mul_sub_back:wwwn #4;{0000}{0000}{0011}{4707}{4559}{7730};
21539     \__fp_fixed_mul_sub_back:wwwn #4;{0000}{0000}{2087}{6756}{9878}{6810};
21540     \__fp_fixed_mul_sub_back:wwwn #4;{0000}{0027}{5573}{1922}{3985}{8907};
21541     \__fp_fixed_mul_sub_back:wwwn #4;{0000}{2480}{1587}{3015}{8730}{1587};
21542     \__fp_fixed_mul_sub_back:wwwn #4;{0013}{8888}{8888}{8888}{8888}{8889};
21543     \__fp_fixed_mul_sub_back:wwwn #4;{0416}{6666}{6666}{6666}{6666}{6667};
21544     \__fp_fixed_mul_sub_back:wwwn #4;{5000}{0000}{0000}{0000}{0000}{0000};
21545     \__fp_fixed_mul_sub_back:wwwn#4;{10000}{0000}{0000}{0000}{0000}{0000};
21546     { \__fp_fixed_continue:wn 0, }
21547   }
21548   { % 1/17!
21549     \__fp_fixed_mul_sub_back:wwwn {0000}{0000}{0000}{0028}{1145}{7254};
21550     #4;{0000}{0000}{0000}{7647}{1637}{3182};
21551     \__fp_fixed_mul_sub_back:wwwn #4;{0000}{0000}{0160}{5904}{3836}{8216};
21552     \__fp_fixed_mul_sub_back:wwwn #4;{0000}{0002}{5052}{1083}{8544}{1719};
21553     \__fp_fixed_mul_sub_back:wwwn #4;{0000}{0275}{5731}{9223}{9858}{9065};
21554     \__fp_fixed_mul_sub_back:wwwn #4;{0001}{9841}{2698}{4126}{9841}{2698};
21555     \__fp_fixed_mul_sub_back:wwwn #4;{0083}{3333}{3333}{3333}{3333}{3333};
21556     \__fp_fixed_mul_sub_back:wwwn #4;{1666}{6666}{6666}{6666}{6666}{6667};
21557     \__fp_fixed_mul_sub_back:wwwn#4;{10000}{0000}{0000}{0000}{0000}{0000};
21558     { \__fp_ep_mul:wwwn 0, } #5,#6;
21559   }
21560   {
21561     \exp_after:wN \__fp_sanitizew
21562     \exp_after:wN #2
21563     \int_value:w \__fp_int_eval:w #1
21564   }

```

```

21565     #2
21566 }

```

(End definition for `_fp_sin_series_o:NNwww` and `_fp_sin_series_aux_o:NNwww`.)

```

\_fp_tan_series_o:NNwww
\_fp_tan_series_aux_o:Nnwww

```

Contrarily to `_fp_sin_series_o:NNwww` which received a conversion auxiliary as `#1`, here, `#1` is 0 for tangent and 2 for cotangent. Consider first the case of the tangent. The octant `#3` starts at 1, which means that it is 1 or 2 for $|x| \in [0, \pi/2]$, it is 3 or 4 for $|x| \in [\pi/2, \pi]$, and so on: the intervals on which $\tan|x| \geq 0$ coincide with those for which $\lfloor (\#3 + 1)/2 \rfloor$ is odd. We also have to take into account the original sign of x to get the sign of the final result; it is straightforward to check that the first `\int_value:w` expansion produces 0 for a positive final result, and 2 otherwise. A similar story holds for $\cot(x)$.

The auxiliary receives the sign, the octant, the square of the (reduced) input, and the (reduced) input (an extended-precision number) as arguments. It then computes the numerator and denominator of

$$\tan(x) \simeq \frac{x(1 - x^2(a_1 - x^2(a_2 - x^2(a_3 - x^2(a_4 - x^2a_5))))))}{1 - x^2(b_1 - x^2(b_2 - x^2(b_3 - x^2(b_4 - x^2b_5)))}.$$

The ratio is computed by `_fp_ep_div:wwwn`, then converted to a floating point number. For octants `#3` (really, quadrants) next to a pole of the functions, the fixed point numerator and denominator are exchanged before computing the ratio. Note that this `\if_int_odd:w` test relies on the fact that the octant is at least 1.

```

21567 \cs_new:Npn \_fp_tan_series_o:NNwww #1#2#3. #4;
21568 {
21569   \_fp_fixed_mul:wwn #4; #4;
21570   {
21571     \exp_after:wN \_fp_tan_series_aux_o:Nnwww
21572     \int_value:w
21573     \if_int_odd:w \_fp_int_eval:w #3 / 2 \_fp_int_eval_end:
21574     \exp_after:wN \reverse_if:N
21575     \fi:
21576     \if_meaning:w #1#2 2 \else: 0 \fi:
21577     {#3}
21578   }
21579 }
21580 \cs_new:Npn \_fp_tan_series_aux_o:Nnwww #1 #2 #3; #4,#5;
21581 {
21582   \_fp_fixed_mul_sub_back:wwwn {0000}{0000}{1527}{3493}{0856}{7059};
21583   #3; {0000}{0159}{6080}{0274}{5257}{6472};
21584   \_fp_fixed_mul_sub_back:wwwn #3; {0002}{4571}{2320}{0157}{2558}{8481};
21585   \_fp_fixed_mul_sub_back:wwwn #3; {0115}{5830}{7533}{5397}{3168}{2147};
21586   \_fp_fixed_mul_sub_back:wwwn #3; {1929}{8245}{6140}{3508}{7719}{2982};
21587   \_fp_fixed_mul_sub_back:wwwn #3; {10000}{0000}{0000}{0000}{0000}{0000};
21588   { \_fp_ep_mul:wwwn 0, } #4,#5;
21589   {
21590     \_fp_fixed_mul_sub_back:wwwn {0000}{0007}{0258}{0681}{9408}{4706};
21591     #3; {0000}{2343}{7175}{1399}{6151}{7670};
21592     \_fp_fixed_mul_sub_back:wwwn #3; {0019}{2638}{4588}{9232}{8861}{3691};
21593     \_fp_fixed_mul_sub_back:wwwn #3; {0536}{6357}{0691}{4344}{6852}{4252};
21594     \_fp_fixed_mul_sub_back:wwwn #3; {5263}{1578}{9473}{6842}{1052}{6315};
21595     \_fp_fixed_mul_sub_back:wwwn #3; {10000}{0000}{0000}{0000}{0000}{0000};

```

```

21596     {
21597         \reverse_if:N \if_int_odd:w
21598         \__fp_int_eval:w (#2 - 1) / 2 \__fp_int_eval_end:
21599         \exp_after:wN \__fp_reverse_args:Nww
21600         \fi:
21601         \__fp_ep_div:wwwn 0,
21602     }
21603 }
21604 {
21605     \exp_after:wN \__fp_sanitizew
21606     \exp_after:wN #1
21607     \int_value:w \__fp_int_eval:w \__fp_ep_to_float_o:wwN
21608 }
21609 #1
21610 }

```

(End definition for `__fp_tan_series_o:NNwww` and `__fp_tan_series_aux_o:Nnwww`.)

34.2 Inverse trigonometric functions

All inverse trigonometric functions (arcsine, arccosine, arctangent, arccotangent, arcsecant, and arcsecant) are based on a function often denoted `atan2`. This function is accessed directly by feeding two arguments to arctangent, and is defined by $\text{atan}(y, x) = \text{atan}(y/x)$ for generic y and x . Its advantages over the conventional arctangent is that it takes values in $[-\pi, \pi]$ rather than $[-\pi/2, \pi/2]$, and that it is better behaved in boundary cases. Other inverse trigonometric functions are expressed in terms of `atan` as

$$\arccos x = \text{atan}(\sqrt{1 - x^2}, x) \quad (5)$$

$$\arcsin x = \text{atan}(x, \sqrt{1 - x^2}) \quad (6)$$

$$\text{asec } x = \text{atan}(\sqrt{x^2 - 1}, 1) \quad (7)$$

$$\text{acsc } x = \text{atan}(1, \sqrt{x^2 - 1}) \quad (8)$$

$$\text{atan } x = \text{atan}(x, 1) \quad (9)$$

$$\text{acot } x = \text{atan}(1, x). \quad (10)$$

Rather than introducing a new function, `atan2`, the arctangent function `atan` is overloaded: it can take one or two arguments. In the comments below, following many texts, we call the first argument y and the second x , because $\text{atan}(y, x) = \text{atan}(y/x)$ is the angular coordinate of the point (x, y) .

As for direct trigonometric functions, the first step in computing $\text{atan}(y, x)$ is argument reduction. The sign of y gives that of the result. We distinguish eight regions where the point $(x, |y|)$ can lie, of angular size roughly $\pi/8$, characterized by their “octant”, between 0 and 7 included. In each region, we compute an arctangent as a Taylor series, then shift this arctangent by the appropriate multiple of $\pi/4$ and sign to get the result. Here is a list of octants, and how we compute the arctangent (we assume $y > 0$; otherwise replace y by $-y$ below):

0 $0 < |y| < 0.41421x$, then $\text{atan } \frac{|y|}{x}$ is given by a nicely convergent Taylor series;

1 $0 < 0.41421x < |y| < x$, then $\text{atan } \frac{|y|}{x} = \frac{\pi}{4} - \text{atan } \frac{x - |y|}{x + |y|}$;

- 2 $0 < 0.41421|y| < x < |y|$, then $\operatorname{atan} \frac{|y|}{x} = \frac{\pi}{4} + \operatorname{atan} \frac{-x+|y|}{x+|y|}$;
- 3 $0 < x < 0.41421|y|$, then $\operatorname{atan} \frac{|y|}{x} = \frac{\pi}{2} - \operatorname{atan} \frac{x}{|y|}$;
- 4 $0 < -x < 0.41421|y|$, then $\operatorname{atan} \frac{|y|}{x} = \frac{\pi}{2} + \operatorname{atan} \frac{-x}{|y|}$;
- 5 $0 < 0.41421|y| < -x < |y|$, then $\operatorname{atan} \frac{|y|}{x} = \frac{3\pi}{4} - \operatorname{atan} \frac{x+|y|}{-x+|y|}$;
- 6 $0 < -0.41421x < |y| < -x$, then $\operatorname{atan} \frac{|y|}{x} = \frac{3\pi}{4} + \operatorname{atan} \frac{-x-|y|}{-x+|y|}$;
- 7 $0 < |y| < -0.41421x$, then $\operatorname{atan} \frac{|y|}{x} = \pi - \operatorname{atan} \frac{|y|}{-x}$.

In the following, we denote by z the ratio among $|\frac{y}{x}|$, $|\frac{x}{y}|$, $|\frac{x+y}{x-y}|$, $|\frac{x-y}{x+y}|$ which appears in the right-hand side above.

34.2.1 Arctangent and arccotangent

`__fp_atan_o:Nw` The parsing step manipulates `atan` and `acot` like `min` and `max`, reading in an array of operands, but also leaves `\use_i:nn` or `\use_ii:nn` depending on whether the result should be given in radians or in degrees. The helper `__fp_parse_function_one_two:nnw` checks that the operand is one or two floating point numbers (not tuples) and leaves its second argument or its tail accordingly (its first argument is used for error messages). More precisely if we are given a single floating point number `__fp_atan_default:w` places `\c_one_fp` (expanded) after it; otherwise `__fp_atan_default:w` is omitted by `__fp_parse_function_one_two:nnw`.

```

21611 \cs_new:Npn \__fp_atan_o:Nw #1
21612 {
21613   \__fp_parse_function_one_two:nnw
21614   { #1 { atan } { atand } }
21615   { \__fp_atan_default:w \__fp_atanii_o:Nww #1 }
21616 }
21617 \cs_new:Npn \__fp_acot_o:Nw #1
21618 {
21619   \__fp_parse_function_one_two:nnw
21620   { #1 { acot } { acotd } }
21621   { \__fp_atan_default:w \__fp_acotii_o:Nww #1 }
21622 }
21623 \cs_new:Npx \__fp_atan_default:w #1#2#3 @ { #1 #2 #3 \c_one_fp @ }

```

(End definition for `__fp_atan_o:Nw`, `__fp_acot_o:Nw`, and `__fp_atan_default:w`.)

`__fp_atanii_o:Nww` If either operand is `nan`, we return it. If both are normal, we call `__fp_atan_normal_o:NNnwNNw`. If both are zero or both infinity, we call `__fp_atan_inf_o:NNNw` with argument 2, leading to a result among $\{\pm\pi/4, \pm3\pi/4\}$ (in degrees, $\{\pm45, \pm135\}$). Otherwise, one is much bigger than the other, and we call `__fp_atan_inf_o:NNNw` with either an argument of 4, leading to the values $\pm\pi/2$ (in degrees, ±90), or 0, leading to $\{\pm0, \pm\pi\}$ (in degrees, $\{\pm0, \pm180\}$). Since `acot(x,y) = atan(y,x)`, `__fp_acotii_o:ww` simply reverses its two arguments.

```

21624 \cs_new:Npn \__fp_atanii_o:Nww
21625   #1 \s__fp \__fp_chk:w #2#3#4; \s__fp \__fp_chk:w #5 #6 @
21626   {
21627     \if_meaning:w 3 #2 \__fp_case_return_i_o:ww \fi:

```

```

21628 \if_meaning:w 3 #5 \__fp_case_return_ii_o:ww \fi:
21629 \if_case:w
21630 \if_meaning:w #2 #5
21631 \if_meaning:w 1 #2 10 \else: 0 \fi:
21632 \else:
21633 \if_int_compare:w #2 > #5 \exp_stop_f: 1 \else: 2 \fi:
21634 \fi:
21635 \exp_stop_f:
21636 \__fp_case_return:nw { \__fp_atan_inf_o:NNNw #1 #3 2 }
21637 \or: \__fp_case_return:nw { \__fp_atan_inf_o:NNNw #1 #3 4 }
21638 \or: \__fp_case_return:nw { \__fp_atan_inf_o:NNNw #1 #3 0 }
21639 \fi:
21640 \__fp_atan_normal_o:NNnwNnw #1
21641 \s__fp \__fp_chk:w #2#3#4;
21642 \s__fp \__fp_chk:w #5 #6
21643 }
21644 \cs_new:Npn \__fp_acotii_o:Nww #1#2; #3;
21645 { \__fp_atanii_o:Nww #1#3; #2; }

```

(End definition for __fp_atanii_o:Nww and __fp_acotii_o:Nww.)

__fp_atan_inf_o:NNNw This auxiliary is called whenever one number is ± 0 or $\pm \infty$ (and neither is NaN). Then the result only depends on the signs, and its value is a multiple of $\pi/4$. We use the same auxiliary as for normal numbers, __fp_atan_combine_o:NwwwwwN, with arguments the final sign #2; the octant #3; $\operatorname{atan} z/z = 1$ as a fixed point number; $z = 0$ as a fixed point number; and $z = 0$ as an extended-precision number. Given the values we provide, $\operatorname{atan} z$ is computed to be 0, and the result is $[#3/2] \cdot \pi/4$ if the sign #5 of x is positive, and $[(7 - #3)/2] \cdot \pi/4$ for negative x , where the divisions are rounded up.

```

21646 \cs_new:Npn \__fp_atan_inf_o:NNNw #1#2#3 \s__fp \__fp_chk:w #4#5#6;
21647 {
21648 \exp_after:wN \__fp_atan_combine_o:NwwwwwN
21649 \exp_after:wN #2
21650 \int_value:w \__fp_int_eval:w
21651 \if_meaning:w 2 #5 7 - \fi: #3 \exp_after:wN ;
21652 \c__fp_one_fixed_tl
21653 {0000}{0000}{0000}{0000}{0000}{0000};
21654 0,{0000}{0000}{0000}{0000}{0000}{0000}; #1
21655 }

```

(End definition for __fp_atan_inf_o:NNNw.)

__fp_atan_normal_o:NNnwNnw Here we simply reorder the floating point data into a pair of signed extended-precision numbers, that is, a sign, an exponent ending with a comma, and a six-block mantissa ending with a semi-colon. This extended precision is required by other inverse trigonometric functions, to compute things like $\operatorname{atan}(x, \sqrt{1 - x^2})$ without intermediate rounding errors.

```

21656 \cs_new_protected:Npn \__fp_atan_normal_o:NNnwNnw
21657 #1 \s__fp \__fp_chk:w 1#2#3#4; \s__fp \__fp_chk:w 1#5#6#7;
21658 {
21659 \__fp_atan_test_o:NwwNwwN
21660 #2 #3, #4{0000}{0000};
21661 #5 #6, #7{0000}{0000}; #1
21662 }

```


(End definition for _fp_atan_normal_o:NNwNnw.)

_fp_atan_test_o:NwwNwwN

This receives: the sign #1 of y , its exponent #2, its 24 digits #3 in groups of 4, and similarly for x . We prepare to call _fp_atan_combine_o:NwwwwwN which expects the sign #1, the octant, the ratio $(\operatorname{atan} z)/z = 1 - \dots$, and the value of z , both as a fixed point number and as an extended-precision floating point number with a mantissa in $[0.01, 1)$. For now, we place #1 as a first argument, and start an integer expression for the octant. The sign of x does not affect z , so we simply leave a contribution to the octant: $\langle \text{octant} \rangle \rightarrow 7 - \langle \text{octant} \rangle$ for negative x . Then we order $|y|$ and $|x|$ in a non-decreasing order: if $|y| > |x|$, insert 3- in the expression for the octant, and swap the two numbers. The finer test with 0.41421 is done by _fp_atan_div:wnwwnw after the operands have been ordered.

```

21663 \cs_new:Npn \_fp_atan_test_o:NwwNwwN #1#2,#3; #4#5,#6;
21664 {
21665   \exp_after:wN \_fp_atan_combine_o:NwwwwwN
21666   \exp_after:wN #1
21667   \int_value:w \_fp_int_eval:w
21668   \if_meaning:w 2 #4
21669     7 - \_fp_int_eval:w
21670   \fi:
21671   \if_int_compare:w
21672     \_fp_ep_compare:www #2,#3; #5,#6; > 0 \exp_stop_f:
21673     3 -
21674   \exp_after:wN \_fp_reverse_args:Nww
21675   \fi:
21676   \_fp_atan_div:wnwwnw #2,#3; #5,#6;
21677 }

```

(End definition for _fp_atan_test_o:NwwNwwN.)

_fp_atan_div:wnwwnw
_fp_atan_near:wwwN
_fp_atan_near_aux:wwN

This receives two positive numbers a and b (equal to $|x|$ and $|y|$ in some order), each as an exponent and 6 blocks of 4 digits, such that $0 < a < b$. If $0.41421b < a$, the two numbers are “near”, hence the point (y, x) that we started with is closer to the diagonals $\{|y| = |x|\}$ than to the axes $\{xy = 0\}$. In that case, the octant is 1 (possibly combined with the 7- and 3- inserted earlier) and we wish to compute $\operatorname{atan} \frac{b-a}{a+b}$. Otherwise, the octant is 0 (again, combined with earlier terms) and we wish to compute $\operatorname{atan} \frac{a}{b}$. In any case, call _fp_atan_auxi:ww followed by z , as a comma-delimited exponent and a fixed point number.

```

21678 \cs_new:Npn \_fp_atan_div:wnwwnw #1,#2#3; #4,#5#6;
21679 {
21680   \if_int_compare:w
21681     \_fp_int_eval:w 41421 * #5 < #2 000
21682     \if_case:w \_fp_int_eval:w #4 - #1 \_fp_int_eval_end:
21683       00 \or: 0 \fi:
21684   \exp_stop_f:
21685   \exp_after:wN \_fp_atan_near:wwwN
21686   \fi:
21687   0
21688   \_fp_ep_div:wwwN #1,{#2}#3; #4,{#5}#6;
21689   \_fp_atan_auxi:ww
21690 }
21691 \cs_new:Npn \_fp_atan_near:wwwN
21692   0 \_fp_ep_div:wwwN #1,#2; #3,

```

```

21693 {
21694     1
21695     \__fp_ep_to_fixed:wwn #1 - #3, #2;
21696     \__fp_atan_near_aux:wwn
21697 }
21698 \cs_new:Npn \__fp_atan_near_aux:wwn #1; #2;
21699 {
21700     \__fp_fixed_add:wwn #1; #2;
21701     { \__fp_fixed_sub:wwn #2; #1; { \__fp_ep_div:wwwn 0, } 0, }
21702 }

```

(End definition for __fp_atan_div:wwwnw, __fp_atan_near:wwwn, and __fp_atan_near_aux:wwn.)

__fp_atan_auxi:ww Convert z from a representation as an exponent and a fixed point number in $[0.01, 1)$ to a
 __fp_atan_auxii:w fixed point number only, then set up the call to __fp_atan_Taylor_loop:www, followed
 by the fixed point representation of z and the old representation.

```

21703 \cs_new:Npn \__fp_atan_auxi:ww #1,#2;
21704 { \__fp_ep_to_fixed:wwn #1,#2; \__fp_atan_auxii:w #1,#2; }
21705 \cs_new:Npn \__fp_atan_auxii:w #1;
21706 {
21707     \__fp_fixed_mul:wwn #1; #1;
21708     {
21709         \__fp_atan_Taylor_loop:www 39 ;
21710         {0000}{0000}{0000}{0000}{0000}{0000} ;
21711     }
21712     ! #1;
21713 }

```

(End definition for __fp_atan_auxi:ww and __fp_atan_auxii:w.)

__fp_atan_Taylor_loop:www We compute the series of $(\operatorname{atan} z)/z$. A typical intermediate stage has $\#1 = 2k - 1$,
 __fp_atan_Taylor_break:w $\#2 = \frac{1}{2k+1} - z^2(\frac{1}{2k+3} - z^2(\dots - z^2\frac{1}{39}))$, and $\#3 = z^2$. To go to the next step $k \rightarrow k - 1$,
 we compute $\frac{1}{2k-1}$, then subtract from it z^2 times $\#2$. The loop stops when $k = 0$: then
 $\#2$ is $(\operatorname{atan} z)/z$, and there is a need to clean up all the unnecessary data, end the integer
 expression computing the octant with a semicolon, and leave the result $\#2$ afterwards.

```

21714 \cs_new:Npn \__fp_atan_Taylor_loop:www #1; #2; #3;
21715 {
21716     \if_int_compare:w #1 = -1 \exp_stop_f:
21717     \__fp_atan_Taylor_break:w
21718     \fi:
21719     \exp_after:wN \__fp_fixed_div_int:wwN \c__fp_one_fixed_tl #1;
21720     \__fp_rrot:www \__fp_fixed_mul_sub_back:wwwn #2; #3;
21721     {
21722         \exp_after:wN \__fp_atan_Taylor_loop:www
21723         \int_value:w \__fp_int_eval:w #1 - 2 ;
21724     }
21725     #3;
21726 }
21727 \cs_new:Npn \__fp_atan_Taylor_break:w
21728 \fi: #1 \__fp_fixed_mul_sub_back:wwwn #2; #3 !
21729 { \fi: ; #2 ; }

```

(End definition for __fp_atan_Taylor_loop:www and __fp_atan_Taylor_break:w.)

`__fp_atan_combine_o:NwwwwN` This receives a $\langle sign \rangle$, an $\langle octant \rangle$, a fixed point value of $(\text{atan } z)/z$, a fixed point number z , and another representation of z , as an $\langle exponent \rangle$ and the fixed point number $10^{-\langle exponent \rangle} z$, followed by either `\use_i:nn` (when working in radians) or `\use_ii:nn` (when working in degrees). The function computes the floating point result

$$\langle sign \rangle \left(\left\lceil \frac{\langle octant \rangle}{2} \right\rceil \frac{\pi}{4} + (-1)^{\langle octant \rangle} \frac{\text{atan } z}{z} \cdot z \right), \quad (11)$$

multiplied by $180/\pi$ if working in degrees, and using in any case the most appropriate representation of z . The floating point result is passed to `__fp_sanitize:Nw`, which checks for overflow or underflow. If the octant is 0, leave the exponent #5 for `__fp_sanitize:Nw`, and multiply #3 = $\frac{\text{atan } z}{z}$ with #6, the adjusted z . Otherwise, multiply #3 = $\frac{\text{atan } z}{z}$ with #4 = z , then compute the appropriate multiple of $\frac{\pi}{4}$ and add or subtract the product #3 · #4. In both cases, convert to a floating point with `__fp_fixed_to_float_o:wN`.

```

21730 \cs_new:Npn \__fp_atan_combine_o:NwwwwN #1 #2; #3; #4; #5,#6; #7
21731 {
21732   \exp_after:wN \__fp_sanitize:Nw
21733   \exp_after:wN #1
21734   \int_value:w \__fp_int_eval:w
21735   \if_meaning:w 0 #2
21736     \exp_after:wN \use_i:nn
21737   \else:
21738     \exp_after:wN \use_ii:nn
21739   \fi:
21740   { #5 \__fp_fixed_mul:wwn #3; #6; }
21741   {
21742     \__fp_fixed_mul:wwn #3; #4;
21743     {
21744       \exp_after:wN \__fp_atan_combine_aux:ww
21745       \int_value:w \__fp_int_eval:w #2 / 2 ; #2;
21746     }
21747   }
21748   { #7 \__fp_fixed_to_float_o:wN \__fp_fixed_to_float_rad_o:wN }
21749   #1
21750 }
21751 \cs_new:Npn \__fp_atan_combine_aux:ww #1; #2;
21752 {
21753   \__fp_fixed_mul_short:wwn
21754   {7853}{9816}{3397}{4483}{0961}{5661};
21755   {#1}{0000}{0000};
21756   {
21757     \if_int_odd:w #2 \exp_stop_f:
21758     \exp_after:wN \__fp_fixed_sub:wwn
21759   \else:
21760     \exp_after:wN \__fp_fixed_add:wwn
21761   \fi:
21762 }
21763 }
```

(End definition for `__fp_atan_combine_o:NwwwwN` and `__fp_atan_combine_aux:ww`.)

34.2.2 Arcsine and arccosine

`__fp_asin_o:w` Again, the first argument provided by `l3fp-parse` is `use_i:nn` if we are to work in radians and `use_ii:nn` for degrees. Then comes a floating point number. The arcsine of ± 0 or NaN is the same floating point number. The arcsine of $\pm\infty$ raises an invalid operation exception. Otherwise, call an auxiliary common with `__fp_acos_o:w`, feeding it information about what function is being performed (for “invalid operation” exceptions).

```

21764 \cs_new:Npn __fp_asin_o:w #1 \s_fp __fp_chk:w #2#3; @
21765 {
21766   \if_case:w #2 \exp_stop_f:
21767     __fp_case_return_same_o:w
21768   \or:
21769     __fp_case_use:nw
21770     { __fp_asin_normal_o:NfwNnnnnw #1 { #1 { asin } { asind } } }
21771   \or:
21772     __fp_case_use:nw
21773     { __fp_invalid_operation_o:fw { #1 { asin } { asind } } }
21774   \else:
21775     __fp_case_return_same_o:w
21776   \fi:
21777   \s_fp __fp_chk:w #2 #3;
21778 }

```

(End definition for `__fp_asin_o:w`.)

`__fp_acos_o:w` The arccosine of ± 0 is $\pi/2$ (in degrees, 90). The arccosine of $\pm\infty$ raises an invalid operation exception. The arccosine of NaN is itself. Otherwise, call an auxiliary common with `__fp_sin_o:w`, informing it that it was called by `acos` or `acosd`, and preparing to swap some arguments down the line.

```

21779 \cs_new:Npn __fp_acos_o:w #1 \s_fp __fp_chk:w #2#3; @
21780 {
21781   \if_case:w #2 \exp_stop_f:
21782     __fp_case_use:nw { __fp_atan_inf_o:NNNw #1 0 4 }
21783   \or:
21784     __fp_case_use:nw
21785     {
21786       __fp_asin_normal_o:NfwNnnnnw #1 { #1 { acos } { acosd } }
21787       __fp_reverse_args:Nww
21788     }
21789   \or:
21790     __fp_case_use:nw
21791     { __fp_invalid_operation_o:fw { #1 { acos } { acosd } } }
21792   \else:
21793     __fp_case_return_same_o:w
21794   \fi:
21795   \s_fp __fp_chk:w #2 #3;
21796 }

```

(End definition for `__fp_acos_o:w`.)

`__fp_asin_normal_o:NfwNnnnnw` If the exponent #5 is at most 0, the operand lies within $(-1, 1)$ and the operation is permitted: call `__fp_asin_auxi_o:NnNw` with the appropriate arguments. If the number is exactly ± 1 (the test works because we know that $\#5 \geq 1$, $\#6\#7 \geq 10000000$, $\#8\#9 \geq 0$,

with equality only for ± 1), we also call `__fp_asin_auxi_o:NnNww`. Otherwise, `__fp_use_i:ww` gets rid of the `asin` auxiliary, and raises instead an invalid operation, because the operand is outside the domain of arcsine or arccosine.

```

21797 \cs_new:Npn \__fp_asin_normal_o:NfwNnnnnw
21798   #1#2#3 \s__fp \__fp_chk:w 1#4#5#6#7#8#9;
21799   {
21800     \if_int_compare:w #5 < 1 \exp_stop_f:
21801       \exp_after:wN \__fp_use_none_until_s:w
21802     \fi:
21803     \if_int_compare:w \__fp_int_eval:w #5 + #6#7 + #8#9 = 1000 0001 ~
21804       \exp_after:wN \__fp_use_none_until_s:w
21805     \fi:
21806     \__fp_use_i:ww
21807     \__fp_invalid_operation_o:fw {#2}
21808     \s__fp \__fp_chk:w 1#4#{#5}{#6}{#7}{#8}{#9};
21809     \__fp_asin_auxi_o:NnNww
21810     #1 {#3} #4 #5,{#6}{#7}{#8}{#9}{0000}{0000};
21811   }

```

(End definition for `__fp_asin_normal_o:NfwNnnnnw`.)

`__fp_asin_auxi_o:NnNww`
`__fp_asin_isqrt:wn`

We compute $x/\sqrt{1-x^2}$. This function is used by `asin` and `acos`, but also by `acsc` and `asec` after inverting the operand, thus it must manipulate extended-precision numbers. First evaluate $1-x^2$ as $(1+x)(1-x)$: this behaves better near $x = 1$. We do the addition/subtraction with fixed point numbers (they are not implemented for extended-precision floats), but go back to extended-precision floats to multiply and compute the inverse square root $1/\sqrt{1-x^2}$. Finally, multiply by the (positive) extended-precision float $|x|$, and feed the (signed) result, and the number $+1$, as arguments to the arctangent function. When computing the arccosine, the arguments $x/\sqrt{1-x^2}$ and $+1$ are swapped by `#2` (`__fp_reverse_args:Nww` in that case) before `__fp_atan_test_o:NwwNwwN` is evaluated. Note that the arctangent function requires normalized arguments, hence the need for `ep_to_ep` and continue after `ep_mul`.

```

21812 \cs_new:Npn \__fp_asin_auxi_o:NnNww #1#2#3#4,#5;
21813   {
21814     \__fp_ep_to_fixed:wwn #4,#5;
21815     \__fp_asin_isqrt:wn
21816     \__fp_ep_mul:wwwwn #4,#5;
21817     \__fp_ep_to_ep:wwN
21818     \__fp_fixed_continue:wn
21819     { #2 \__fp_atan_test_o:NwwNwwN #3 }
21820     0 1,{1000}{0000}{0000}{0000}{0000}{0000}; #1
21821   }
21822 \cs_new:Npn \__fp_asin_isqrt:wn #1;
21823   {
21824     \exp_after:wN \__fp_fixed_sub:wwn \c__fp_one_fixed_tl #1;
21825     {
21826       \__fp_fixed_add_one:wn #1;
21827       \__fp_fixed_continue:wn { \__fp_ep_mul:wwwwn 0, } 0,
21828     }
21829     \__fp_ep_isqrt:wwn
21830   }

```

(End definition for `__fp_asin_auxi_o:NnNww` and `__fp_asin_isqrt:wn`.)

34.2.3 Arccosecant and arcsecant

`__fp_acsc_o:w` Cases are mostly labelled by #2, except when #2 is 2: then we use #3#2, which is 02 = 2 when the number is $+\infty$ and 22 when the number is $-\infty$. The arccosecant of ± 0 raises an invalid operation exception. The arccosecant of $\pm\infty$ is ± 0 with the same sign. The arcosecant of NaN is itself. Otherwise, `__fp_acsc_normal_o:NfwNnw` does some more tests, keeping the function name (acsc or acscd) as an argument for invalid operation exceptions.

```

21831 \cs_new:Npn \__fp_acsc_o:w #1 \s__fp \__fp_chk:w #2#3#4; @
21832 {
21833   \if_case:w \if_meaning:w 2 #2 #3 \fi: #2 \exp_stop_f:
21834     \__fp_case_use:nw
21835     { \__fp_invalid_operation_o:fw { #1 { acsc } { acscd } } }
21836   \or: \__fp_case_use:nw
21837     { \__fp_acsc_normal_o:NfwNnw #1 { #1 { acsc } { acscd } } }
21838   \or: \__fp_case_return_o:Nw \c_zero_fp
21839   \or: \__fp_case_return_same_o:w
21840   \else: \__fp_case_return_o:Nw \c_minus_zero_fp
21841   \fi:
21842   \s__fp \__fp_chk:w #2 #3 #4;
21843 }

```

(End definition for `__fp_acsc_o:w`.)

`__fp_asec_o:w` The arcsecant of ± 0 raises an invalid operation exception. The arcsecant of $\pm\infty$ is $\pi/2$ (in degrees, 90). The arcosecant of NaN is itself. Otherwise, do some more tests, keeping the function name asec (or asecd) as an argument for invalid operation exceptions, and a `__fp_reverse_args:Nww` following precisely that appearing in `__fp_acos_o:w`.

```

21844 \cs_new:Npn \__fp_asec_o:w #1 \s__fp \__fp_chk:w #2#3; @
21845 {
21846   \if_case:w #2 \exp_stop_f:
21847     \__fp_case_use:nw
21848     { \__fp_invalid_operation_o:fw { #1 { asec } { asecd } } }
21849   \or:
21850     \__fp_case_use:nw
21851     {
21852       \__fp_acsc_normal_o:NfwNnw #1 { #1 { asec } { asecd } }
21853       \__fp_reverse_args:Nww
21854     }
21855   \or: \__fp_case_use:nw { \__fp_atan_inf_o:NNNw #1 0 4 }
21856   \else: \__fp_case_return_same_o:w
21857   \fi:
21858   \s__fp \__fp_chk:w #2 #3;
21859 }

```

(End definition for `__fp_asec_o:w`.)

`__fp_acsc_normal_o:NfwNnw` If the exponent is non-positive, the operand is less than 1 in absolute value, which is always an invalid operation: complain. Otherwise, compute the inverse of the operand, and feed it to `__fp_asin_auxi_o:NnNww` (with all the appropriate arguments). This computes what we want thanks to $\text{acsc}(x) = \text{asin}(1/x)$ and $\text{asec}(x) = \text{acos}(1/x)$.

```

21860 \cs_new:Npn \__fp_acsc_normal_o:NfwNnw #1#2#3 \s__fp \__fp_chk:w 1#4#5#6;
21861 {
21862   \int_compare:nNnTF {#5} < 1

```

```

21863     {
21864         \__fp_invalid_operation_o:fw {#2}
21865         \s__fp \__fp_chk:w 1#4{#5}#6;
21866     }
21867     {
21868         \__fp_ep_div:wwwwn
21869         1,{1000}{0000}{0000}{0000}{0000}{0000};
21870         #5,#6{0000}{0000};
21871     { \__fp_asin_auxi_o:NnNww #1 {#3} #4 }
21872     }
21873 }

```

(End definition for __fp_acsc_normal_o:NfwNnw.)

```

21874 \end{package}

```

35 13fp-convert implementation

```

21875 \begin{package}

```

```

21876 \set{fp}

```

35.1 Dealing with tuples

The first argument is for instance __fp_to_t1_dispatch:w, which converts any floating point object to the appropriate representation. We loop through all items, putting ,~ between all of them and making sure to remove the leading ,~.

```

21877 \cs_new:Npn \__fp_tuple_convert:Nw #1 \s__fp_tuple \__fp_tuple_chk:w #2 ;
21878 {
21879     \int_case:nnF { \__fp_array_count:n {#2} }
21880     {
21881         { 0 } { ( ) }
21882         { 1 } { \__fp_tuple_convert_end:w @ { #1 #2 , } }
21883     }
21884     {
21885         \__fp_tuple_convert_loop:nNw { } #1
21886         #2 { ? \__fp_tuple_convert_end:w } ;
21887         @ { \use_none:nn }
21888     }
21889 }
21890 \cs_new:Npn \__fp_tuple_convert_loop:nNw #1#2#3#4; #5 @ #6
21891 {
21892     \use_none:n #3
21893     \exp_args:Nf \__fp_tuple_convert_loop:nNw { #2 #3#4 ; } #2 #5
21894     @ { #6 , ~ #1 }
21895 }
21896 \cs_new:Npn \__fp_tuple_convert_end:w #1 @ #2
21897 { \exp_after:wN ( \exp:w \exp_end_continue_f:w #2 ) }

```

(End definition for __fp_tuple_convert:Nw, __fp_tuple_convert_loop:nNw, and __fp_tuple_convert_end:w.)

35.2 Trimming trailing zeros

`__fp_trim_zeros:w` If #1 ends with a 0, the loop auxiliary takes that zero as an end-delimiter for its first argument, and the second argument is the same loop auxiliary. Once the last trailing zero is reached, the second argument is the dot auxiliary, which removes a trailing dot if any. We then clean-up with the end auxiliary, keeping only the number.

```
21898 \cs_new:Npn __fp_trim_zeros:w #1 ;
21899 {
21900   __fp_trim_zeros_loop:w #1
21901   ; __fp_trim_zeros_loop:w 0; __fp_trim_zeros_dot:w .; \s__fp_stop
21902 }
21903 \cs_new:Npn __fp_trim_zeros_loop:w #1 0; #2 { #2 #1 ; #2 }
21904 \cs_new:Npn __fp_trim_zeros_dot:w #1 .; { __fp_trim_zeros_end:w #1 ; }
21905 \cs_new:Npn __fp_trim_zeros_end:w #1 ; #2 \s__fp_stop { #1 }
```

(End definition for `__fp_trim_zeros:w` and others.)

35.3 Scientific notation

`\fp_to_scientific:N` The three public functions evaluate their argument, then pass it to `__fp_to_scientific_dispatch:w`.

```
\fp_to_scientific:c
\fp_to_scientific:n
21906 \cs_new:Npn \fp_to_scientific:N #1
21907 { \exp_after:wN __fp_to_scientific_dispatch:w #1 }
21908 \cs_generate_variant:Nn \fp_to_scientific:N { c }
21909 \cs_new:Npn \fp_to_scientific:n
21910 {
21911   \exp_after:wN __fp_to_scientific_dispatch:w
21912   \exp:w \exp_end_continue_f:w __fp_parse:n
21913 }
```

(End definition for `\fp_to_scientific:N` and `\fp_to_scientific:n`. These functions are documented on page 206.)

```
\__fp_to_scientific_dispatch:w
\__fp_to_scientific_recover:w
__fp_tuple_to_scientific:w
21914 \cs_new:Npn __fp_to_scientific_dispatch:w #1
21915 {
21916   __fp_change_func_type:NNN
21917   #1 __fp_to_scientific:w __fp_to_scientific_recover:w
21918   #1
21919 }
21920 \cs_new:Npn __fp_to_scientific_recover:w #1 #2 ;
21921 {
21922   __fp_error:nffn { fp-unknown-type } { \tl_to_str:n { #2 ; } } { } { }
21923   nan
21924 }
21925 \cs_new:Npn __fp_tuple_to_scientific:w
21926 { __fp_tuple_convert:Nw __fp_to_scientific_dispatch:w }
```

(End definition for `__fp_to_scientific_dispatch:w`, `__fp_to_scientific_recover:w`, and `__fp_tuple_to_scientific:w`.)

`__fp_to_scientific:w` Expressing an internal floating point number in scientific notation is quite easy: no rounding, and the format is very well defined. First cater for the sign: negative numbers (#2 = 2) start with -; we then only need to care about positive numbers and `nan`. Then

filter the special cases: ± 0 are represented as 0; infinities are converted to a number slightly larger than the largest after an “invalid_operation” exception; `nan` is represented as 0 after an “invalid_operation” exception. In the normal case, decrement the exponent and unbrace the 4 brace groups, then in a second step grab the first digit (previously hidden in braces) to order the various parts correctly.

```

21927 \cs_new:Npn \__fp_to_scientific:w \s__fp \__fp_chk:w #1#2
21928 {
21929   \if_meaning:w 2 #2 \exp_after:wN - \exp:w \exp_end_continue_f:w \fi:
21930   \if_case:w #1 \exp_stop_f:
21931     \__fp_case_return:nw { 0.000000000000000e0 }
21932   \or: \exp_after:wN \__fp_to_scientific_normal:wnnnnn
21933   \or:
21934     \__fp_case_use:nw
21935     {
21936       \__fp_invalid_operation:nnw
21937       { \fp_to_scientific:N \c__fp_overflowing_fp }
21938       { fp_to_scientific }
21939     }
21940   \or:
21941     \__fp_case_use:nw
21942     {
21943       \__fp_invalid_operation:nnw
21944       { \fp_to_scientific:N \c_zero_fp }
21945       { fp_to_scientific }
21946     }
21947   \fi:
21948   \s__fp \__fp_chk:w #1 #2
21949 }
21950 \cs_new:Npn \__fp_to_scientific_normal:wnnnnn
21951   \s__fp \__fp_chk:w 1 #1 #2 #3#4#5#6 ;
21952   {
21953     \exp_after:wN \__fp_to_scientific_normal:wNw
21954     \exp_after:wN e
21955     \int_value:w \__fp_int_eval:w #2 - 1
21956     ; #3 #4 #5 #6 ;
21957   }
21958 \cs_new:Npn \__fp_to_scientific_normal:wNw #1 ; #2#3;
21959   { #2.#3 #1 }

```

(End definition for `__fp_to_scientific:w`, `__fp_to_scientific_normal:wnnnnn`, and `__fp_to_scientific_normal:wNw`.)

35.4 Decimal representation

`\fp_to_decimal:N` All three public variants are based on the same `__fp_to_decimal_dispatch:w` after evaluating their argument to an internal floating point.

```

\fp_to_decimal:c
\fp_to_decimal:n
21960 \cs_new:Npn \fp_to_decimal:N #1
21961   { \exp_after:wN \__fp_to_decimal_dispatch:w #1 }
21962 \cs_generate_variant:Nn \fp_to_decimal:N { c }
21963 \cs_new:Npn \fp_to_decimal:n
21964   {
21965     \exp_after:wN \__fp_to_decimal_dispatch:w
21966     \exp:w \exp_end_continue_f:w \__fp_parse:n
21967   }

```

(End definition for `\fp_to_decimal:N` and `\fp_to_decimal:n`. These functions are documented on page 206.)

`__fp_to_decimal_dispatch:w`
`__fp_to_decimal_recover:w`
`__fp_tuple_to_decimal:w`

We allow tuples.

```

21968 \cs_new:Npn \__fp_to_decimal_dispatch:w #1
21969 {
21970   \__fp_change_func_type:NNN
21971   #1 \__fp_to_decimal:w \__fp_to_decimal_recover:w
21972   #1
21973 }
21974 \cs_new:Npn \__fp_to_decimal_recover:w #1 #2 ;
21975 {
21976   \__fp_error:nffn { fp-unknown-type } { \tl_to_str:n { #2 ; } } { } { }
21977   nan
21978 }
21979 \cs_new:Npn \__fp_tuple_to_decimal:w
21980 { \__fp_tuple_convert:Nw \__fp_to_decimal_dispatch:w }
```

(End definition for `__fp_to_decimal_dispatch:w`, `__fp_to_decimal_recover:w`, and `__fp_tuple_to_decimal:w`.)

`__fp_to_decimal:w`
`_fp_to_decimal_normal:wnnnnn`
`__fp_to_decimal_large:Nnnw`
`__fp_to_decimal_huge:wnnnn`

The structure is similar to `__fp_to_scientific:w`. Insert `-` for negative numbers. Zero gives 0, $\pm\infty$ and NaN yield an “invalid operation” exception; note that $\pm\infty$ produces a very large output, which we don’t expand now since it most likely won’t be needed. Normal numbers with an exponent in the range [1, 15] have that number of digits before the decimal separator: “decimate” them, and remove leading zeros with `\int_value:w`, then trim trailing zeros and dot. Normal numbers with an exponent 16 or larger have no decimal separator, we only need to add trailing zeros. When the exponent is non-positive, the result should be 0.<zeros><digits>, trimmed.

```

21981 \cs_new:Npn \__fp_to_decimal:w \s__fp \__fp_chk:w #1#2
21982 {
21983   \if_meaning:w 2 #2 \exp_after:wN - \exp:w \exp_end_continue_f:w \fi:
21984   \if_case:w #1 \exp_stop_f:
21985     \__fp_case_return:nw { 0 }
21986   \or: \exp_after:wN \__fp_to_decimal_normal:wnnnnn
21987   \or:
21988     \__fp_case_use:nw
21989     {
21990       \__fp_invalid_operation:nnw
21991       { \fp_to_decimal:N \c__fp_overflowing_fp }
21992       { fp_to_decimal }
21993     }
21994   \or:
21995     \__fp_case_use:nw
21996     {
21997       \__fp_invalid_operation:nnw
21998       { 0 }
21999       { fp_to_decimal }
22000     }
22001   \fi:
22002   \s__fp \__fp_chk:w #1 #2
22003 }
22004 \cs_new:Npn \__fp_to_decimal_normal:wnnnnn
22005 \s__fp \__fp_chk:w 1 #1 #2 #3#4#5#6 ;
```

```

22006 {
22007   \int_compare:nNnTF {#2} > 0
22008   {
22009     \int_compare:nNnTF {#2} < \c__fp_prec_int
22010     {
22011       \__fp_decimate:nNnnnn { \c__fp_prec_int - #2 }
22012       \__fp_to_decimal_large:Nnnw
22013     }
22014     {
22015       \exp_after:wN \exp_after:wN
22016       \exp_after:wN \__fp_to_decimal_huge:wnnnn
22017       \prg_replicate:nn { #2 - \c__fp_prec_int } { 0 } ;
22018     }
22019     {#3} {#4} {#5} {#6}
22020   }
22021   {
22022     \exp_after:wN \__fp_trim_zeros:w
22023     \exp_after:wN 0
22024     \exp_after:wN .
22025     \exp:w \exp_end_continue_f:w \prg_replicate:nn { - #2 } { 0 }
22026     #3#4#5#6 ;
22027   }
22028 }
22029 \cs_new:Npn \__fp_to_decimal_large:Nnnw #1#2#3#4;
22030 {
22031   \exp_after:wN \__fp_trim_zeros:w \int_value:w
22032   \if_int_compare:w #2 > 0 \exp_stop_f:
22033   #2
22034   \fi:
22035   \exp_stop_f:
22036   #3.#4 ;
22037 }
22038 \cs_new:Npn \__fp_to_decimal_huge:wnnnn #1; #2#3#4#5 { #2#3#4#5 #1 }

```

(End definition for `__fp_to_decimal:w` and others.)

35.5 Token list representation

`\fp_to_tl:N` These three public functions evaluate their argument, then pass it to `__fp_to_tl_dispatch:w`.
`\fp_to_tl:c` dispatch:w.
`\fp_to_tl:n`

```

22039 \cs_new:Npn \fp_to_tl:N #1 { \exp_after:wN \__fp_to_tl_dispatch:w #1 }
22040 \cs_generate_variant:Nn \fp_to_tl:N { c }
22041 \cs_new:Npn \fp_to_tl:n
22042 {
22043   \exp_after:wN \__fp_to_tl_dispatch:w
22044   \exp:w \exp_end_continue_f:w \__fp_parse:n
22045 }

```

(End definition for `\fp_to_tl:N` and `\fp_to_tl:n`. These functions are documented on page 207.)

`__fp_to_tl_dispatch:w` We allow tuples.
`__fp_to_tl_recover:w` 22046 \cs_new:Npn __fp_to_tl_dispatch:w #1
`__fp_tuple_to_tl:w` 22047 { __fp_change_func_type:NNN #1 __fp_to_tl:w __fp_to_tl_recover:w #1 }
22048 \cs_new:Npn __fp_to_tl_recover:w #1 #2 ;

```

22049 {
22050     \__fp_error:nffn { fp-unknown-type } { \tl_to_str:n { #2 ; } } { } { }
22051     nan
22052 }
22053 \cs_new:Npn \__fp_tuple_to_tl:w
22054 { \__fp_tuple_convert:Nw \__fp_to_tl_dispatch:w }

```

(End definition for __fp_to_tl_dispatch:w, __fp_to_tl_recover:w, and __fp_tuple_to_tl:w.)

__fp_to_tl:w A structure similar to __fp_to_scientific_dispatch:w and __fp_to_decimal_dispatch:w, but without the “invalid operation” exception. First filter special cases. We express normal numbers in decimal notation if the exponent is in the range $[-2, 16]$, and otherwise use scientific notation.

```

22055 \cs_new:Npn \__fp_to_tl:w \s__fp \__fp_chk:w #1#2
22056 {
22057     \if_meaning:w 2 #2 \exp_after:wN - \exp:w \exp_end_continue_f:w \fi:
22058     \if_case:w #1 \exp_stop_f:
22059         \__fp_case_return:nw { 0 }
22060     \or: \exp_after:wN \__fp_to_tl_normal:nnnnn
22061     \or: \__fp_case_return:nw { inf }
22062     \else: \__fp_case_return:nw { nan }
22063     \fi:
22064 }
22065 \cs_new:Npn \__fp_to_tl_normal:nnnnn #1
22066 {
22067     \int_compare:nTF
22068     { -2 <= #1 <= \c__fp_prec_int }
22069     { \__fp_to_decimal_normal:wnnnnnn }
22070     { \__fp_to_tl_scientific:wnnnnnn }
22071     \s__fp \__fp_chk:w 1 0 {#1}
22072 }
22073 \cs_new:Npn \__fp_to_tl_scientific:wnnnnnn
22074 \s__fp \__fp_chk:w 1 #1 #2 #3#4#5#6 ;
22075 {
22076     \exp_after:wN \__fp_to_tl_scientific:wNw
22077     \exp_after:wN e
22078     \int_value:w \__fp_int_eval:w #2 - 1
22079     ; #3 #4 #5 #6 ;
22080 }
22081 \cs_new:Npn \__fp_to_tl_scientific:wNw #1 ; #2#3;
22082 { \__fp_trim_zeros:w #2.#3 ; #1 }

```

(End definition for __fp_to_tl:w and others.)

35.6 Formatting

This is not implemented yet, as it is not yet clear what a correct interface would be, for this kind of structured conversion from a floating point (or other types of variables) to a string. Ideas welcome.

35.7 Convert to dimension or integer

`\fp_to_dim:N` All three public variants are based on the same `__fp_to_dim_dispatch:w` after evaluating their argument to an internal floating point. We only allow floating point numbers, not tuples.

```
\__fp_to_dim_dispatch:w 22083 \cs_new:Npn \fp_to_dim:N #1
\__fp_to_dim_recover:w 22084 { \exp_after:wN \__fp_to_dim_dispatch:w #1 }
\__fp_to_dim:w 22085 \cs_generate_variant:Nn \fp_to_dim:N { c }
22086 \cs_new:Npn \fp_to_dim:n
22087 {
22088   \exp_after:wN \__fp_to_dim_dispatch:w
22089   \exp:w \exp_end_continue_f:w \__fp_parse:n
22090 }
22091 \cs_new:Npn \__fp_to_dim_dispatch:w #1#2 ;
22092 {
22093   \__fp_change_func_type:NNN #1 \__fp_to_dim:w \__fp_to_dim_recover:w
22094   #1 #2 ;
22095 }
22096 \cs_new:Npn \__fp_to_dim_recover:w #1
22097 { \__fp_invalid_operation:nnw { Opt } { fp_to_dim } }
22098 \cs_new:Npn \__fp_to_dim:w #1 ; { \__fp_to_decimal:w #1 ; pt }
```

(End definition for `\fp_to_dim:N` and others. These functions are documented on page 206.)

`\fp_to_int:N` For the most part identical to `\fp_to_dim:N` but without `pt`, and where `__fp_to_int:w` does more work. To convert to an integer, first round to 0 places (to the nearest integer), then express the result as a decimal number: the definition of `__fp_to_decimal_dispatch:w` is such that there are no trailing dot nor zero.

```
\__fp_to_int_dispatch:w 22099 \cs_new:Npn \fp_to_int:N #1 { \exp_after:wN \__fp_to_int_dispatch:w #1 }
\__fp_to_int_recover:w 22100 \cs_generate_variant:Nn \fp_to_int:N { c }
22101 \cs_new:Npn \fp_to_int:n
22102 {
22103   \exp_after:wN \__fp_to_int_dispatch:w
22104   \exp:w \exp_end_continue_f:w \__fp_parse:n
22105 }
22106 \cs_new:Npn \__fp_to_int_dispatch:w #1#2 ;
22107 {
22108   \__fp_change_func_type:NNN #1 \__fp_to_int:w \__fp_to_int_recover:w
22109   #1 #2 ;
22110 }
22111 \cs_new:Npn \__fp_to_int_recover:w #1
22112 { \__fp_invalid_operation:nnw { 0 } { fp_to_int } }
22113 \cs_new:Npn \__fp_to_int:w #1;
22114 {
22115   \exp_after:wN \__fp_to_decimal:w \exp:w \exp_end_continue_f:w
22116   \__fp_round:Nwn \__fp_round_to_nearest:NNN #1; { 0 }
22117 }
```

(End definition for `\fp_to_int:N` and others. These functions are documented on page 206.)

35.8 Convert from a dimension

`\dim_to_fp:n` The dimension expression (which can in fact be a glue expression) is evaluated, converted to a number (*i.e.*, expressed in scaled points), then multiplied by $2^{-16} =$

```
\__fp_from_dim_test:ww
\__fp_from_dim:wNw
\__fp_from_dim:wNNnnnnnn
\__fp_from_dim:wnnnnwNw
```

0.0000152587890625 to give a value expressed in points. The auxiliary `__fp_mul_npos_o:Nww` expects the desired *final sign* and two floating point operands (of the form `\s__fp ... ;`) as arguments. This set of functions is also used to convert dimension registers to floating points while parsing expressions: in this context there is an additional exponent, which is the first argument of `__fp_from_dim_test:ww`, and is combined with the exponent -4 of 2^{-16} . There is also a need to expand afterwards: this is performed by `__fp_mul_npos_o:Nww`, and cancelled by `\prg_do_nothing:` here.

```

22118 \cs_new:Npn \dim_to_fp:n #1
22119 {
22120   \exp_after:wN \__fp_from_dim_test:ww
22121   \exp_after:wN 0
22122   \exp_after:wN ,
22123   \int_value:w \tex_glueexpr:D #1 ;
22124 }
22125 \cs_new:Npn \__fp_from_dim_test:ww #1, #2
22126 {
22127   \if_meaning:w 0 #2
22128     \__fp_case_return:nw { \exp_after:wN \c_zero_fp }
22129   \else:
22130     \exp_after:wN \__fp_from_dim:wNw
22131     \int_value:w \__fp_int_eval:w #1 - 4
22132     \if_meaning:w - #2
22133       \exp_after:wN , \exp_after:wN 2 \int_value:w
22134     \else:
22135       \exp_after:wN , \exp_after:wN 0 \int_value:w #2
22136     \fi:
22137   \fi:
22138 }
22139 \cs_new:Npn \__fp_from_dim:wNw #1,#2#3;
22140 {
22141   \__fp_pack_twice_four:wNNNNNNNN \__fp_from_dim:wNNnnnnnn ;
22142   #3 000 0000 00 {10}987654321; #2 {#1}
22143 }
22144 \cs_new:Npn \__fp_from_dim:wNNnnnnnn #1; #2#3#4#5#6#7#8#9
22145 { \__fp_from_dim:wnnnnwNn #1 {#2#300} {0000} ; }
22146 \cs_new:Npn \__fp_from_dim:wnnnnwNn #1; #2#3#4#5#6; #7#8
22147 {
22148   \__fp_mul_npos_o:Nww #7
22149   \s__fp \__fp_chk:w 1 #7 {#5} #1 ;
22150   \s__fp \__fp_chk:w 1 0 {#8} {1525} {8789} {0625} {0000} ;
22151   \prg_do_nothing:
22152 }

```

(End definition for `\dim_to_fp:n` and others. This function is documented on page 179.)

35.9 Use and eval

`\fp_use:N` Those public functions are simple copies of the decimal conversions.
`\fp_use:c`
`\fp_eval:n`

```

22153 \cs_new_eq:NN \fp_use:N \fp_to_decimal:N
22154 \cs_generate_variant:Nn \fp_use:N { c }
22155 \cs_new_eq:NN \fp_eval:n \fp_to_decimal:n

```

(End definition for `\fp_use:N` and `\fp_eval:n`. These functions are documented on page 207.)

\fp_sign:n Trivial but useful. See the implementation of `\fp_add:Nn` for an explanation of why to use `__fp_parse:n`, namely, for better error reporting.

```
22156 \cs_new:Npn \fp_sign:n #1
22157 { \fp_to_decimal:n { sign \__fp_parse:n {#1} } }
```

(End definition for `\fp_sign:n`. This function is documented on page 206.)

\fp_abs:n Trivial but useful. See the implementation of `\fp_add:Nn` for an explanation of why to use `__fp_parse:n`, namely, for better error reporting.

```
22158 \cs_new:Npn \fp_abs:n #1
22159 { \fp_to_decimal:n { abs \__fp_parse:n {#1} } }
```

(End definition for `\fp_abs:n`. This function is documented on page 221.)

\fp_max:nn Similar to `\fp_abs:n`, for consistency with `\int_max:nn`, etc.

\fp_min:nn

```
22160 \cs_new:Npn \fp_max:nn #1#2
22161 { \fp_to_decimal:n { max ( \__fp_parse:n {#1} , \__fp_parse:n {#2} ) } }
22162 \cs_new:Npn \fp_min:nn #1#2
22163 { \fp_to_decimal:n { min ( \__fp_parse:n {#1} , \__fp_parse:n {#2} ) } }
```

(End definition for `\fp_max:nn` and `\fp_min:nn`. These functions are documented on page 221.)

35.10 Convert an array of floating points to a comma list

`__fp_array_to_clist:n` Converts an array of floating point numbers to a comma-list. If speed here ends up irrelevant, we can simplify the code for the auxiliary to become

```
\cs_new:Npn \__fp_array_to_clist_loop:Nw #1#2;
{
  \use_none:n #1
  { , ~ } \fp_to_tl:n { #1 #2 ; }
  \__fp_array_to_clist_loop:Nw
}
```

The `\use_ii:nn` function is expanded after `__fp_expand:n` is done, and it removes `,~` from the start of the representation.

```
22164 \cs_new:Npn \__fp_array_to_clist:n #1
22165 {
22166   \tl_if_empty:nF {#1}
22167   {
22168     \exp_last_unbraced:Ne \use_ii:nn
22169     {
22170       \__fp_array_to_clist_loop:Nw #1 { ? \prg_break: } ;
22171       \prg_break_point:
22172     }
22173   }
22174 }
22175 \cs_new:Npn \__fp_array_to_clist_loop:Nw #1#2;
22176 {
22177   \use_none:n #1
22178   , ~
22179   \exp_not:f { \__fp_to_tl_dispatch:w #1 #2 ; }
22180   \__fp_array_to_clist_loop:Nw
22181 }
```

(End definition for `_fp_array_to_clist:n` and `_fp_array_to_clist_loop:Nw`.)

22182 `\</package>`

36 13fp-random Implementation

22183 `*package>`

22184 `<@@=fp>`

`_fp_parse_word_rand:N` Those functions may receive a variable number of arguments. We won't use the argument `?`.
`_fp_parse_word_randint:N`

22185 `\cs_new:Npn _fp_parse_word_rand:N`
22186 `{ _fp_parse_function:NNN _fp_rand_o:Nw ? }`
22187 `\cs_new:Npn _fp_parse_word_randint:N`
22188 `{ _fp_parse_function:NNN _fp_randint_o:Nw ? }`

(End definition for `_fp_parse_word_rand:N` and `_fp_parse_word_randint:N`.)

36.1 Engine support

Most engines provide random numbers, but not all. We write the test twice simply in order to write the `false` branch first.

22189 `\sys_if_rand_exist:F`
22190 `{`
22191 `_kernel_msg_new:nnn { kernel } { fp-no-random }`
22192 `{ Random-numbers-unavailable-for~#1 }`
22193 `\cs_new:Npn _fp_rand_o:Nw ? #1 @`
22194 `{`
22195 `_kernel_msg_expandable_error:nnn { kernel } { fp-no-random }`
22196 `{ fp-rand }`
22197 `\exp_after:wN \c_nan_fp`
22198 `}`
22199 `\cs_new_eq:NN _fp_randint_o:Nw _fp_rand_o:Nw`
22200 `\cs_new:Npn \int_rand:nn #1#2`
22201 `{`
22202 `_kernel_msg_expandable_error:nnn { kernel } { fp-no-random }`
22203 `{ \int_rand:nn {#1} {#2} }`
22204 `\int_eval:n {#1}`
22205 `}`
22206 `\cs_new:Npn \int_rand:n #1`
22207 `{`
22208 `_kernel_msg_expandable_error:nnn { kernel } { fp-no-random }`
22209 `{ \int_rand:n {#1} }`
22210 `1`
22211 `}`
22212 `}`
22213 `\sys_if_rand_exist:T`
22214 `{`

Obviously, every word “random” below means “pseudo-random”, as we have no access to entropy (except a very unreliable source of entropy: the time it takes to run some code).

The primitive random number generator (RNG) is provided as `\tex_uniformdeviate:D`. Under the hood, it maintains an array of 55 28-bit numbers, updated with a linear recursion relation (similar to Fibonacci numbers) modulo 2^{28} . When `\tex_uniformdeviate:D` $\langle integer \rangle$ is called (for brevity denote by N the $\langle integer \rangle$), the next 28-bit number is read from the array, scaled by $N/2^{28}$, and rounded. To prevent 0 and N from appearing half as often as other numbers, they are both mapped to the result 0.

This process means that `\tex_uniformdeviate:D` only gives a uniform distribution from 0 to $N-1$ if N is a divisor of 2^{28} , so we will mostly call the RNG with such power of 2 arguments. If N does not divide 2^{28} , then the relative non-uniformity (difference between probabilities of getting different numbers) is about $N/2^{28}$. This implies that detecting deviation from $1/N$ of the probability of a fixed value X requires about $2^{56}/N$ random trials. But collective patterns can reduce this to about $2^{56}/N^2$. For instance with $N = 3 \times 2^k$, the modulo 3 repartition of such random numbers is biased with a non-uniformity about $2^k/2^{28}$ (which is much worse than the circa $3/2^{28}$ non-uniformity from taking directly $N = 3$). This is detectable after about $2^{56}/2^{2k} = 9 \cdot 2^{56}/N^2$ random numbers. For $k = 15$, $N = 98304$, this means roughly 2^{26} calls to the RNG (experimentally this takes at the very least 16 seconds on a 2 giga-hertz processor). While this bias is not quite problematic, it is uncomfortably close to being so, and it becomes worse as N is increased. In our code, we shall thus combine several results from the RNG.

The RNG has three types of unexpected correlations. First, everything is linear modulo 2^{28} , hence the lowest k bits of the random numbers only depend on the lowest k bits of the seed (and of course the number of times the RNG was called since setting the seed). The recommended way to get a number from 0 to $N-1$ is thus to scale the raw 28-bit integer, as the engine's RNG does. We will go further and in fact typically we discard some of the lowest bits.

Second, suppose that we call the RNG with the same argument N to get a set of K integers in $[0, N-1]$ (throwing away repeats), and suppose that $N > K^3$ and $K > 55$. The recursion used to construct more 28-bit numbers from previous ones is linear: $x_n = x_{n-55} - x_{n-24}$ or $x_n = x_{n-55} - x_{n-24} + 2^{28}$. After rescaling and rounding we find that the result $N_n \in [0, N-1]$ is among $N_{n-55} - N_{n-24} + \{-1, 0, 1\}$ modulo N (a more detailed analysis shows that 0 appears with frequency close to $3/4$). The resulting set thus has more triplets (a, b, c) than expected obeying $a = b + c$ modulo N . Namely it will have of order $(K-55) \times 3/4$ such triplets, when one would expect $K^3/(6N)$. This starts to be detectable around $N = 2^{18} > 55^3$ (earlier if one keeps track of positions too, but this is more subtle than it looks because the array of 28-bit integers is read backwards by the engine). Hopefully the correlation is subtle enough to not affect realistic documents so we do not specifically mitigate against this. Since we typically use two calls to the RNG per `\int_rand:nn` we would need to investigate linear relations between the x_{2n} on the one hand and between the x_{2n+1} on the other hand. Such relations will have more complicated coefficients than ± 1 , which alleviates the issue.

Third, consider successive batches of 165 calls to the RNG (with argument 2^{28} or with argument 2 for instance), then most batches have more odd than even numbers. Note that this does not mean that there are more odd than even numbers overall. Similar issues are discussed in Knuth's TAOCP volume 2 near exercise 3.3.2-31. We do not have any mitigation strategy for this.

Ideally, our algorithm should be:

- Uniform. The result should be as uniform as possible assuming that the RNG's underlying 28-bit integers are uniform.

- Uncorrelated. The result should not have detectable correlations between different seeds, similar to the lowest-bit ones mentioned earlier.
- Quick. The algorithm should be fast in \TeX , so no “bit twiddling”, but “digit twiddling” is ok.
- Simple. The behaviour must be documentable precisely.
- Predictable. The number of calls to the RNG should be the same for any `\int_rand:nn`, because then the algorithm can be modified later without changing the result of other uses of the RNG.
- Robust. It should work even for `\int_rand:nn { - \c_max_int } { \c_max_int }` where the range is not representable as an integer. In fact, we also provide later a floating-point `randint` whose range can go all the way up to $2 \times 10^{16} - 1$ possible values.

Some of these requirements conflict. For instance, uniformity cannot be achieved with a fixed number of calls to the RNG.

Denote by `random(N)` one call to `\tex_uniformdeviate:D` with argument N , and by `ediv(p, q)` the ε - \TeX rounding division giving $\lfloor p/q + 1/2 \rfloor$. Denote by $\langle \min \rangle$, $\langle \max \rangle$ and $R = \langle \max \rangle - \langle \min \rangle + 1$ the arguments of `\int_min:nn` and the number of possible outcomes. Note that $R \in [1, 2^{32} - 1]$ cannot necessarily be represented as an integer (however, $R - 2^{31}$ can). Our strategy is to get two 28-bit integers X and Y from the RNG, split each into 14-bit integers, as $X = X_1 \times 2^{14} + X_0$ and $Y = Y_1 \times 2^{14} + Y_0$ then return essentially $\langle \min \rangle + \lfloor R(X_1 \times 2^{-14} + Y_1 \times 2^{-28} + Y_0 \times 2^{-42} + X_0 \times 2^{-56}) \rfloor$. For small R the X_0 term has a tiny effect so we ignore it and we can compute $R \times Y/2^{28}$ much more directly by `random(R)`.

- If $R \leq 2^{17} - 1$ then return `ediv(R random(2^{14}) + random(R) + 2^{13} , 2^{14}) - 1 + $\langle \min \rangle$` . The shifts by 2^{13} and -1 convert ε - \TeX division to truncated division. The bound on R ensures that the number obtained after the shift is less than `\c_max_int`. The non-uniformity is at most of order $2^{17}/2^{42} = 2^{-25}$.
- Split $R = R_2 \times 2^{28} + R_1 \times 2^{14} + R_0$, where $R_2 \in [0, 15]$. Compute $\langle \min \rangle + R_2 X_1 2^{14} + (R_2 Y_1 + R_1 X_1) + \text{ediv}(R_2 Y_0 + R_1 Y_1 + R_0 X_1 + \text{ediv}(R_2 X_0 + R_0 Y_1 + \text{ediv}((2^{14} R_1 + R_0)(2^{14} Y_0 + X_0), 2^{28}), 2^{14}), 2^{14})$ then map a result of $\langle \max \rangle + 1$ to $\langle \min \rangle$. Writing each `ediv` in terms of truncated division with a shift, and using $\lfloor (p + \lfloor r/s \rfloor)/q \rfloor = \lfloor (ps + r)/(sq) \rfloor$, what we compute is equal to $\lfloor \langle \text{exact} \rangle + 2^{-29} + 2^{-15} + 2^{-1} \rfloor$ with $\langle \text{exact} \rangle = \langle \min \rangle + R \times 0.X_1 Y_1 Y_0 X_0$. Given we map $\langle \max \rangle + 1$ to $\langle \min \rangle$, the shift has no effect on uniformity. The non-uniformity is bounded by $R/2^{56} < 2^{-24}$. It may be possible to speed up the code by dropping tiny terms such as $R_0 X_0$, but the analysis of non-uniformity proves too difficult.

To avoid the overflow when the computation yields $\langle \max \rangle + 1$ with $\langle \max \rangle = 2^{31} - 1$ (note that R is then arbitrary), we compute the result in two pieces. Compute $\langle \text{first} \rangle = \langle \min \rangle + R_2 X_1 2^{14}$ if $R_2 < 8$ or $\langle \min \rangle + 8 X_1 2^{14} + (R_2 - 8) X_1 2^{14}$ if $R_2 \geq 8$, the expressions being chosen to avoid overflow. Compute $\langle \text{second} \rangle = R_2 Y_1 + R_1 X_1 + \text{ediv}(\dots)$, at most $R_2 2^{14} + R_1 2^{14} + R_0 \leq 2^{28} + 15 \times 2^{14} - 1$, not at risk of overflowing. We have $\langle \text{first} \rangle + \langle \text{second} \rangle = \langle \max \rangle + 1 = \langle \min \rangle + R$ if and only if $\langle \text{second} \rangle = R 2^{14} + R_0 + R_2 2^{14}$ and $2^{14} R_2 X_1 = 2^{28} R_2 - 2^{14} R_2$ (namely $R_2 = 0$ or $X_1 = 2^{14} - 1$). In that case, return $\langle \min \rangle$, otherwise return $\langle \text{first} \rangle + \langle \text{second} \rangle$, which is safe because it is at most $\langle \max \rangle$. Note that the decision of what to return

does not need $\langle first \rangle$ explicitly so we don't actually compute it, just put it in an integer expression in which $\langle second \rangle$ is eventually added (or not).

- To get a floating point number in $[0, 1)$ just call the $R = 10000 \leq 2^{17} - 1$ procedure above to produce four blocks of four digits.
- To get an integer floating point number in a range (whose size can be up to $2 \times 10^{16} - 1$), work with fixed-point numbers: get six times four digits to build a fixed point number, multiply by R and add $\langle min \rangle$. This requires some care because l3fp-extended only supports non-negative numbers.

`\c__kernel_randint_max_int` Constant equal to $2^{17} - 1$, the maximal size of a range that `\int_range:nn` can do with its “simple” algorithm.

```
22215 \int_const:Nn \c__kernel_randint_max_int { 131071 }
```

(End definition for `\c__kernel_randint_max_int`.)

`__kernel_randint:n` Used in an integer expression, `__kernel_randint:n {R}` gives a random number $1 + \lfloor (R \text{random}(2^{14}) + \text{random}(R)) / 2^{14} \rfloor$ that is in $[1, R]$. Previous code was computing $\lfloor p / 2^{14} \rfloor$ as `ediv(p - 2^{13}, 2^{14})` but that wrongly gives -1 for $p = 0$.

```
22216 \cs_new:Npn \__kernel_randint:n #1
22217 {
22218   (#1 * \tex_uniformdeviate:D 16384
22219   + \tex_uniformdeviate:D #1 + 8192 ) / 16384
22220 }
```

(End definition for `__kernel_randint:n`.)

`__fp_rand_myriads:n` Used as `__fp_rand_myriads:n {XXX}` with one letter X (specifically) per block of four digit we want; it expands to `;` followed by the requested number of brace groups, each containing four (pseudo-random) digits. Digits are produced as a random number in $[10000, 19999]$ for the usual reason of preserving leading zeros.

```
22221 \cs_new:Npn \__fp_rand_myriads:n #1
22222 { \__fp_rand_myriads_loop:w #1 \prg_break: X \prg_break_point: ; }
22223 \cs_new:Npn \__fp_rand_myriads_loop:w #1 X
22224 {
22225   #1
22226   \exp_after:wN \__fp_rand_myriads_get:w
22227   \int_value:w \__fp_int_eval:w 9999 +
22228   \__kernel_randint:n { 10000 }
22229   \__fp_rand_myriads_loop:w
22230 }
22231 \cs_new:Npn \__fp_rand_myriads_get:w 1 #1 ; { ; {#1} }
```

(End definition for `__fp_rand_myriads:n`, `__fp_rand_myriads_loop:w`, and `__fp_rand_myriads_get:w`.)

36.2 Random floating point

`__fp_rand_o:Nw` First we check that `random` was called without argument. Then get four blocks of four digits and convert that fixed point number to a floating point number (this correctly sets the exponent). This has a minor bug: if all of the random numbers are zero then the result is correctly 0 but it raises the underflow flag; it should not do that.

`__fp_rand_o:w`

```

22232 \cs_new:Npn \__fp_rand_o:Nw ? #1 @
22233 {
22234   \tl_if_empty:nTF {#1}
22235   {
22236     \exp_after:wN \__fp_rand_o:w
22237     \exp:w \exp_end_continue_f:w
22238     \__fp_rand_myriads:n { XXXX } { 0000 } { 0000 } ; 0
22239   }
22240   {
22241     \__kernel_msg_expandable_error:nnnnn
22242     { kernel } { fp-num-args } { rand() } { 0 } { 0 }
22243     \exp_after:wN \c_nan_fp
22244   }
22245 }
22246 \cs_new:Npn \__fp_rand_o:w ;
22247 {
22248   \exp_after:wN \__fp_sanitizew:Nw
22249   \exp_after:wN 0
22250   \int_value:w \__fp_int_eval:w \c_zero_int
22251   \__fp_fixed_to_float_o:wN
22252 }

```

(End definition for `__fp_rand_o:Nw` and `__fp_rand_o:w`.)

36.3 Random integer

`__fp_randint_o:Nw` Enforce that there is one argument (then add first argument 1) or two arguments. Call `__fp_randint_badarg:w` on each; this function inserts `1 \exp_stop_f:` to end the `\if_case:w` statement if either the argument is not an integer or if its absolute value is $\geq 10^{16}$. Also bail out if `__fp_compare_back:ww` yields 1, meaning that the bounds are not in the right order. Otherwise an auxiliary converts each argument times 10^{-16} (hence the shift in exponent) to a 24-digit fixed point number (see `l3fp-extended`). Then compute the number of choices, $\langle max \rangle + 1 - \langle min \rangle$. Create a random 24-digit fixed-point number with `__fp_rand_myriads:n`, then use a fused multiply-add instruction to multiply the number of choices to that random number and add it to $\langle min \rangle$. Then truncate to 16 digits (namely select the integer part of 10^{16} times the result) before converting back to a floating point number (`__fp_sanitizew:Nw` takes care of zero). To avoid issues with negative numbers, add 1 to all fixed point numbers (namely 10^{16} to the integers they represent), except of course when it is time to convert back to a float.

```

22253 \cs_new:Npn \__fp_randint_o:Nw ?
22254 {
22255   \__fp_parse_function_one_two:nnw
22256   { randint }
22257   { \__fp_randint_default:w \__fp_randint_o:w }
22258 }
22259 \cs_new:Npn \__fp_randint_default:w #1 { \exp_after:wN #1 \c_one_fp }
22260 \cs_new:Npn \__fp_randint_badarg:w \s__fp \__fp_chk:w #1#2#3;

```

```

22261 {
22262     \__fp_int:wTF \s__fp \__fp_chk:w #1#2#3;
22263     {
22264         \if_meaning:w 1 #1
22265         \if_int_compare:w
22266             \__fp_use_i_until_s:nw #3 ; > \c__fp_prec_int
22267             1 \exp_stop_f:
22268         \fi:
22269         \fi:
22270     }
22271     { 1 \exp_stop_f: }
22272 }
22273 \cs_new:Npn \__fp_randint_o:w #1; #2; @
22274 {
22275     \if_case:w
22276         \__fp_randint_badarg:w #1;
22277         \__fp_randint_badarg:w #2;
22278         \if:w 1 \__fp_compare_back:ww #2; #1; 1 \exp_stop_f: \fi:
22279         0 \exp_stop_f:
22280         \__fp_randint_auxi_o:ww #1; #2;
22281     \or:
22282         \__fp_invalid_operation_tl_o:ff
22283         { randint } { \__fp_array_to_clist:n { #1; #2; } }
22284     \exp:w
22285     \fi:
22286     \exp_after:wN \exp_end:
22287 }
22288 \cs_new:Npn \__fp_randint_auxi_o:ww #1 ; #2 ; #3 \exp_end:
22289 {
22290     \fi:
22291     \__fp_randint_auxii:wn #2 ;
22292     { \__fp_randint_auxii:wn #1 ; \__fp_randint_auxiii_o:ww }
22293 }
22294 \cs_new:Npn \__fp_randint_auxii:wn \s__fp \__fp_chk:w #1#2#3#4 ;
22295 {
22296     \if_meaning:w 0 #1
22297     \exp_after:wN \use_i:nn
22298     \else:
22299     \exp_after:wN \use_ii:nn
22300     \fi:
22301     { \exp_after:wN \__fp_fixed_continue:wn \c__fp_one_fixed_tl }
22302     {
22303         \exp_after:wN \__fp_ep_to_fixed:wwn
22304         \int_value:w \__fp_int_eval:w
22305         #3 - \c__fp_prec_int , #4 {0000} {0000} ;
22306         {
22307             \if_meaning:w 0 #2
22308             \exp_after:wN \use_i:nnnn
22309             \exp_after:wN \__fp_fixed_add_one:wN
22310             \fi:
22311             \exp_after:wN \__fp_fixed_sub:wwn \c__fp_one_fixed_tl
22312         }
22313         \__fp_fixed_continue:wn
22314     }

```

```

22315     }
22316 \cs_new:Npn \__fp_randint_auxiii_o:ww #1 ; #2 ;
22317 {
22318     \__fp_fixed_add:wwn #2 ;
22319     {0000} {0000} {0000} {0001} {0000} {0000} ;
22320     \__fp_fixed_sub:wwn #1 ;
22321     {
22322         \exp_after:wN \use_i:nn
22323         \exp_after:wN \__fp_fixed_mul_add:wwwn
22324         \exp:w \exp_end_continue_f:w \__fp_rand_myriads:n { XXXXXX } ;
22325     }
22326     #1 ;
22327     \__fp_randint_auxiv_o:ww
22328     #2 ;
22329     \__fp_randint_auxv_o:w #1 ; @
22330 }
22331 \cs_new:Npn \__fp_randint_auxiv_o:ww #1#2#3#4#5 ; #6#7#8#9
22332 {
22333     \if_int_compare:w
22334         \if_int_compare:w #1#2 > #6#7 \exp_stop_f: 1 \else:
22335         \if_int_compare:w #1#2 < #6#7 \exp_stop_f: - \fi: \fi:
22336         #3#4 > #8#9 \exp_stop_f:
22337         \__fp_use_i_until_s:nw
22338         \fi:
22339         \__fp_randint_auxv_o:w {#1}{#2}{#3}{#4}#5
22340 }
22341 \cs_new:Npn \__fp_randint_auxv_o:w #1#2#3#4#5 ; #6 @
22342 {
22343     \exp_after:wN \__fp_sanitize:Nw
22344     \int_value:w
22345     \if_int_compare:w #1 < 10000 \exp_stop_f:
22346         2
22347     \else:
22348         0
22349         \exp_after:wN \exp_after:wN
22350         \exp_after:wN \__fp_reverse_args:Nww
22351     \fi:
22352     \exp_after:wN \__fp_fixed_sub:wwn \c__fp_one_fixed_tl
22353     {#1} {#2} {#3} {#4} {0000} {0000} ;
22354     {
22355         \exp_after:wN \exp_stop_f:
22356         \int_value:w \__fp_int_eval:w \c__fp_prec_int
22357         \__fp_fixed_to_float_o:wN
22358     }
22359     0
22360     \exp:w \exp_after:wN \exp_end:
22361 }

```

(End definition for __fp_randint_o:Nw and others.)

\int_rand:nn Evaluate the argument and filter out the case where the lower bound #1 is more than
__fp_randint:ww the upper bound #2. Then determine whether the range is narrower than **\c__kernel_-randint_max_int**; #2-#1 may overflow for very large positive #2 and negative #1. If the range is narrow, call **__kernel_randint:n {⟨choices⟩}** where ⟨choices⟩ is the number

of possible outcomes. If the range is wide, use somewhat slower code.

```

22362 \cs_new:Npn \int_rand:nn #1#2
22363 {
22364   \int_eval:n
22365   {
22366     \exp_after:wN \__fp_randint:ww
22367     \int_value:w \int_eval:n {#1} \exp_after:wN ;
22368     \int_value:w \int_eval:n {#2} ;
22369   }
22370 }
22371 \cs_new:Npn \__fp_randint:ww #1; #2;
22372 {
22373   \if_int_compare:w #1 > #2 \exp_stop_f:
22374   \__kernel_msg_expandable_error:nnnn
22375   { kernel } { randint-backward-range } {#1} {#2}
22376   \__fp_randint:ww #2; #1;
22377   \else:
22378     \if_int_compare:w \__fp_int_eval:w #2
22379     \if_int_compare:w #1 > \c_zero_int
22380     - #1 < \__fp_int_eval:w
22381     \else:
22382       < \__fp_int_eval:w #1 +
22383       \fi:
22384       \c_kernel_randint_max_int
22385       \__fp_int_eval_end:
22386       \__kernel_randint:n
22387       { \__fp_int_eval:w #2 - #1 + 1 \__fp_int_eval_end: }
22388       - 1 + #1
22389     \else:
22390       \__kernel_randint:nn {#1} {#2}
22391     \fi:
22392   \fi:
22393 }

```

(End definition for `\int_rand:nn` and `__fp_randint:ww`. This function is documented on page 99.)

Any $n \in [-2^{31} + 1, 2^{31} - 1]$ is uniquely written as $2^{14}n_1 + n_2$ with $n_1 \in [-2^{17}, 2^{17} - 1]$ and $n_2 \in [0, 2^{14} - 1]$. Calling `__fp_randint_split_o:Nw n` ; gives n_1 ; n_2 ; and expands the next token once. We do this for two random numbers and apply `__fp_randint_split_o:Nw` twice to fully decompose the range R . One subtlety is that we compute $R - 2^{31} = \langle \max \rangle - \langle \min \rangle - (2^{31} - 1) \in [-2^{31} + 1, 2^{31} - 1]$ rather than R to avoid overflow.

Then we have `__fp_randint_wide_aux:w` $\langle X_1 \rangle; \langle X_0 \rangle; \langle Y_1 \rangle; \langle Y_0 \rangle; \langle R_2 \rangle; \langle R_1 \rangle; \langle R_0 \rangle; .$ and we apply the algorithm described earlier.

```

22394 \cs_new:Npn \__kernel_randint:nn #1#2
22395 {
22396   #1
22397   \exp_after:wN \__fp_randint_wide_aux:w
22398   \int_value:w
22399   \exp_after:wN \__fp_randint_split_o:Nw
22400   \tex_uniformdeviate:D 268435456 ;
22401   \int_value:w
22402   \exp_after:wN \__fp_randint_split_o:Nw
22403   \tex_uniformdeviate:D 268435456 ;
22404   \int_value:w

```

```

22405         \exp_after:wN \__fp_randint_split_o:Nw
22406         \int_value:w \__fp_int_eval:w 131072 +
22407         \exp_after:wN \__fp_randint_split_o:Nw
22408         \int_value:w
22409         \__kernel_int_add:nnn {#2} { -#1 } { -\c_max_int } ;
22410     .
22411 }
22412 \cs_new:Npn \__fp_randint_split_o:Nw #1#2 ;
22413 {
22414     \if_meaning:w 0 #1
22415     0 \exp_after:wN ; \int_value:w 0
22416     \else:
22417         \exp_after:wN \__fp_randint_split_aux:w
22418         \int_value:w \__fp_int_eval:w (#1#2 - 8192) / 16384 ;
22419         + #1#2
22420     \fi:
22421     \exp_after:wN ;
22422 }
22423 \cs_new:Npn \__fp_randint_split_aux:w #1 ;
22424 {
22425     #1 \exp_after:wN ;
22426     \int_value:w \__fp_int_eval:w - #1 * 16384
22427 }
22428 \cs_new:Npn \__fp_randint_wide_aux:w #1;#2; #3;#4; #5;#6;#7; .
22429 {
22430     \exp_after:wN \__fp_randint_wide_auxii:w
22431     \int_value:w \__fp_int_eval:w #5 * #3 + #6 * #1 +
22432     (#5 * #4 + #6 * #3 + #7 * #1 +
22433     (#5 * #2 + #7 * #3 +
22434     (16384 * #6 + #7) * (16384 * #4 + #2) / 268435456) / 16384
22435     ) / 16384 \exp_after:wN ;
22436     \int_value:w \__fp_int_eval:w (#5 + #6) * 16384 + #7 ;
22437     #1 ; #5 ;
22438 }
22439 \cs_new:Npn \__fp_randint_wide_auxii:w #1; #2; #3; #4;
22440 {
22441     \if_int_odd:w 0
22442     \if_int_compare:w #1 = #2 \else: \exp_stop_f: \fi:
22443     \if_int_compare:w #4 = \c_zero_int 1 \fi:
22444     \if_int_compare:w #3 = 16383 ~ 1 \fi:
22445     \exp_stop_f:
22446     \exp_after:wN \prg_break:
22447     \fi:
22448     \if_int_compare:w #4 < 8 \exp_stop_f:
22449     + #4 * #3 * 16384
22450     \else:
22451     + 8 * #3 * 16384 + (#4 - 8) * #3 * 16384
22452     \fi:
22453     + #1
22454     \prg_break_point:
22455 }

```

(End definition for __kernel_randint:nn and others.)

\int_rand:n Similar to \int_rand:nn, but needs fewer checks.

```

\__fp_randint:n 22456 \cs_new:Npn \int_rand:n #1
22457 {
22458   \int_eval:n
22459   { \exp_args:Nf \__fp_randint:n { \int_eval:n {#1} } }
22460 }
22461 \cs_new:Npn \__fp_randint:n #1
22462 {
22463   \if_int_compare:w #1 < 1 \exp_stop_f:
22464   \__kernel_msg_expandable_error:nnnn
22465   { kernel } { randint-backward-range } { 1 } {#1}
22466   \__fp_randint:ww #1; 1;
22467   \else:
22468   \if_int_compare:w #1 > \c__kernel_randint_max_int
22469   \__kernel_randint:nn { 1 } {#1}
22470   \else:
22471   \__kernel_randint:n {#1}
22472   \fi:
22473   \fi:
22474 }

```

(End definition for \int_rand:n and __fp_randint:n. This function is documented on page 99.)

End the initial conditional that ensures these commands are only defined in engines that support random numbers.

```

22475 }
22476 \</package>

```

37 l3fparray implementation

```

22477 \*package>
22478 \@@=fp>

```

In analogy to l3intarray it would make sense to have <@@=fparray>, but we need direct access to __fp_parse:n from l3fp-parse, and a few other (less crucial) internals of the l3fp family.

37.1 Allocating arrays

There are somewhat more than $(2^{31} - 1)^2$ floating point numbers so we store each floating point number as three entries in integer arrays. To avoid having to multiply indices by three or to add 1 etc, a floating point array is just a token list consisting of three tokens: integer arrays of the same size.

\g__fp_array_int Used to generate unique names for the three integer arrays.

```

22479 \int_new:N \g__fp_array_int

```

(End definition for \g__fp_array_int.)

\l__fp_array_loop_int Used to loop in __fp_array_gzero:N.

```

22480 \int_new:N \l__fp_array_loop_int

```

(End definition for \l__fp_array_loop_int.)

\fpararray_new:Nn Build a three-token token list, then define all three tokens to be integer arrays of the same size. No need to initialize the data: the integer arrays start with zeros, and three zeros denote precisely `\c_zero_fp`, as we want.

\fpararray_new:cn

__fp_array_new:nNNN

```

22481 \cs_new_protected:Npn \fpararray_new:Nn #1#2
22482 {
22483   \tl_new:N #1
22484   \prg_replicate:nn { 3 }
22485   {
22486     \int_gincr:N \g__fp_array_int
22487     \exp_args:NNc \tl_gput_right:Nn #1
22488     { g__fp_array_ \__fp_int_to_roman:w \g__fp_array_int _intarray }
22489   }
22490   \exp_last_unbraced:Nfo \__fp_array_new:nNNNN
22491   { \int_eval:n {#2} } #1 #1
22492 }
22493 \cs_generate_variant:Nn \fpararray_new:Nn { c }
22494 \cs_new_protected:Npn \__fp_array_new:nNNNN #1#2#3#4#5
22495 {
22496   \int_compare:nNnTF {#1} < 0
22497   {
22498     \__kernel_msg_error:nnn { kernel } { negative-array-size } {#1}
22499     \cs_undefine:N #1
22500     \int_gsub:Nn \g__fp_array_int { 3 }
22501   }
22502   {
22503     \intarray_new:Nn #2 {#1}
22504     \intarray_new:Nn #3 {#1}
22505     \intarray_new:Nn #4 {#1}
22506   }
22507 }

```

(End definition for `\fpararray_new:Nn` and `__fp_array_new:nNNN`. This function is documented on page 224.)

\fpararray_count:N Size of any of the intarrays, here we pick the third.

\fpararray_count:c

```

22508 \cs_new:Npn \fpararray_count:N #1
22509 {
22510   \exp_after:wN \use_i:nnn
22511   \exp_after:wN \intarray_count:N #1
22512 }
22513 \cs_generate_variant:Nn \fpararray_count:N { c }

```

(End definition for `\fpararray_count:N`. This function is documented on page 224.)

37.2 Array items

__fp_array_bounds:NNnTF See the `l3intarray` analogue: only names change. The functions `\fpararray_gset:Nnn` and `\fpararray_item:Nn` share bounds checking. The T branch is used if #3 is within bounds of the array #2.

__fp_array_bounds_error:NNn

```

22514 \cs_new:Npn \__fp_array_bounds:NNnTF #1#2#3#4#5
22515 {
22516   \if_int_compare:w 1 > #3 \exp_stop_f:
22517   \__fp_array_bounds_error:NNn #1 #2 {#3}
22518   #5

```

```

22519     \else:
22520         \if_int_compare:w #3 > \fpararray_count:N #2 \exp_stop_f:
22521             \__fp_array_bounds_error:NNn #1 #2 {#3}
22522             #5
22523         \else:
22524             #4
22525         \fi:
22526     \fi:
22527 }
22528 \cs_new:Npn \__fp_array_bounds_error:NNn #1#2#3
22529 {
22530     #1 { kernel } { out-of-bounds }
22531     { \token_to_str:N #2 } {#3} { \fpararray_count:N #2 }
22532 }

```

(End definition for __fp_array_bounds:NNnTF and __fp_array_bounds_error:NNn.)

\fpararray_gset:Nnn

Evaluate, then store exponent in one intarray, sign and 8 digits of mantissa in the next, and 8 trailing digits in the last.

\fpararray_gset:cnn

__fp_array_gset:NNNNnw
__fp_array_gset:w
__fp_array_gset_recover:Nw
__fp_array_gset_special:nnNNN
__fp_array_gset_normal:w

```

22533 \cs_new_protected:Npn \fpararray_gset:Nnn #1#2#3
22534 {
22535     \exp_after:wN \exp_after:wN
22536     \exp_after:wN \__fp_array_gset:NNNNnw
22537     \exp_after:wN #1
22538     \exp_after:wN #1
22539     \int_value:w \int_eval:n {#2} \exp_after:wN ;
22540     \exp:w \exp_end_continue_f:w \__fp_parse:n {#3}
22541 }
22542 \cs_generate_variant:Nn \fpararray_gset:Nnn { c }
22543 \cs_new_protected:Npn \__fp_array_gset:NNNNnw #1#2#3#4#5 ; #6 ;
22544 {
22545     \__fp_array_bounds:NNnTF \__kernel_msg_error:nnxxx #4 {#5}
22546     {
22547         \exp_after:wN \__fp_change_func_type:NNN
22548         \__fp_use_i_until_s:nw #6 ;
22549         \__fp_array_gset:w
22550         \__fp_array_gset_recover:Nw
22551         #6 ; {#5} #1 #2 #3
22552     }
22553     { }
22554 }
22555 \cs_new_protected:Npn \__fp_array_gset_recover:Nw #1#2 ;
22556 {
22557     \__fp_error:nffn { fp-unknown-type } { \tl_to_str:n { #2 ; } } { } { } { }
22558     \exp_after:wN #1 \c_nan_fp
22559 }
22560 \cs_new_protected:Npn \__fp_array_gset:w \s__fp \__fp_chk:w #1#2
22561 {
22562     \if_case:w #1 \exp_stop_f:
22563         \__fp_case_return:nw { \__fp_array_gset_special:nnNNN {#2} }
22564     \or: \exp_after:wN \__fp_array_gset_normal:w
22565     \or: \__fp_case_return:nw { \__fp_array_gset_special:nnNNN { #2 3 } }
22566     \or: \__fp_case_return:nw { \__fp_array_gset_special:nnNNN { 1 } }
22567     \fi:

```

```

22568     \s__fp \__fp_chk:w #1 #2
22569   }
22570 \cs_new_protected:Npn \__fp_array_gset_normal:w
22571   \s__fp \__fp_chk:w 1 #1 #2 #3#4#5 ; #6#7#8#9
22572   {
22573     \__kernel_intarray_gset:Nnn #7 {#6} {#2}
22574     \__kernel_intarray_gset:Nnn #8 {#6}
22575     { \if_meaning:w 2 #1 3 \else: 1 \fi: #3#4 }
22576     \__kernel_intarray_gset:Nnn #9 {#6} { 1 \use:nn #5 }
22577   }
22578 \cs_new_protected:Npn \__fp_array_gset_special:nnNNN #1#2#3#4#5
22579   {
22580     \__kernel_intarray_gset:Nnn #3 {#2} {#1}
22581     \__kernel_intarray_gset:Nnn #4 {#2} {0}
22582     \__kernel_intarray_gset:Nnn #5 {#2} {0}
22583   }

```

(End definition for \fpararray_gset:Nnn and others. This function is documented on page 224.)

\fpararray_gzero:N

\fpararray_gzero:c

```

22584 \cs_new_protected:Npn \fpararray_gzero:N #1
22585   {
22586     \int_zero:N \l__fp_array_loop_int
22587     \prg_replicate:nn { \fpararray_count:N #1 }
22588     {
22589       \int_incr:N \l__fp_array_loop_int
22590       \exp_after:wN \__fp_array_gset_special:nnNNN
22591       \exp_after:wN 0
22592       \exp_after:wN \l__fp_array_loop_int
22593       #1
22594     }
22595   }
22596 \cs_generate_variant:Nn \fpararray_gzero:N { c }

```

(End definition for \fpararray_gzero:N. This function is documented on page 224.)

\fpararray_item:Nn

\fpararray_item:cn

\fpararray_item_to_tl:Nn

\fpararray_item_to_tl:cn

__fp_array_item:NwN

__fp_array_item:NNNnN

__fp_array_item:N

__fp_array_item:w

__fp_array_item_special:w

__fp_array_item_normal:w

```

22597 \cs_new:Npn \fpararray_item:Nn #1#2
22598   {
22599     \exp_after:wN \__fp_array_item:NwN
22600     \exp_after:wN #1
22601     \int_value:w \int_eval:n {#2} ;
22602     \__fp_to_decimal:w
22603   }
22604 \cs_generate_variant:Nn \fpararray_item:Nn { c }
22605 \cs_new:Npn \fpararray_item_to_tl:Nn #1#2
22606   {
22607     \exp_after:wN \__fp_array_item:NwN
22608     \exp_after:wN #1
22609     \int_value:w \int_eval:n {#2} ;
22610     \__fp_to_tl:w
22611   }
22612 \cs_generate_variant:Nn \fpararray_item_to_tl:Nn { c }
22613 \cs_new:Npn \__fp_array_item:NwN #1#2 ; #3
22614   {

```

```

22615     \_fp_array_bounds:NNnTF \_kernel_msg_expandable_error:nnfff #1 {#2}
22616     { \exp_after:wN \_fp_array_item:NNNnN #1 {#2} #3 }
22617     { \exp_after:wN #3 \c_nan_fp }
22618   }
22619 \cs_new:Npn \_fp_array_item:NNNnN #1#2#3#4
22620 {
22621   \exp_after:wN \_fp_array_item:N
22622   \int_value:w \_kernel_intarray_item:Nn #2 {#4} \exp_after:wN ;
22623   \int_value:w \_kernel_intarray_item:Nn #3 {#4} \exp_after:wN ;
22624   \int_value:w \_kernel_intarray_item:Nn #1 {#4} ;
22625 }
22626 \cs_new:Npn \_fp_array_item:N #1
22627 {
22628   \if_meaning:w 0 #1 \exp_after:wN \_fp_array_item_special:w \fi:
22629   \_fp_array_item:w #1
22630 }
22631 \cs_new:Npn \_fp_array_item:w #1 #2#3#4#5 #6 ; 1 #7 ;
22632 {
22633   \exp_after:wN \_fp_array_item_normal:w
22634   \int_value:w \if_meaning:w #1 1 0 \else: 2 \fi: \exp_stop_f:
22635   #7 ; {#2#3#4#5} {#6} ;
22636 }
22637 \cs_new:Npn \_fp_array_item_special:w #1 ; #2 ; #3 ; #4
22638 {
22639   \exp_after:wN #4
22640   \exp:w \exp_end_continue_f:w
22641   \if_case:w #3 \exp_stop_f:
22642     \exp_after:wN \c_zero_fp
22643   \or: \exp_after:wN \c_nan_fp
22644   \or: \exp_after:wN \c_minus_zero_fp
22645   \or: \exp_after:wN \c_inf_fp
22646   \else: \exp_after:wN \c_minus_inf_fp
22647   \fi:
22648 }
22649 \cs_new:Npn \_fp_array_item_normal:w #1 #2#3#4#5 #6 ; #7 ; #8 ; #9
22650 { #9 \s_fp \_fp_chk:w 1 #1 {#8} #7 {#2#3#4#5} {#6} ; }

```

(End definition for `\fparray_item:Nn` and others. These functions are documented on page 224.)

```

22651 \endpackage

```

38 l3cctab implementation

```

22652 \*package>
22653 \@@=cctab>

```

As LuaTeX offers engine support for category code tables, and this is entirely lacking from the other engines, we need two complementary approaches. (Some future XeTeX may add support, at which point the conditionals below would be different.)

38.1 Variables

`\g__cctab_stack_seq` List of catcode tables saved by nested `\cctab_begin:N`, to restore catcodes at the matching `\cctab_end:.` When popped from the `\g__cctab_stack_seq` the table numbers are stored in `\g__cctab_unused_seq` for later reuse.

```

22654 \seq_new:N \g__cctab_stack_seq
22655 \seq_new:N \g__cctab_unused_seq

```

(End definition for `\g__cctab_stack_seq` and `\g__cctab_unused_seq`.)

`\g__cctab_group_seq` A stack to store the group level when a catcode table started.

```

22656 \seq_new:N \g__cctab_group_seq

```

(End definition for `\g__cctab_group_seq`.)

`\g__cctab_allocate_int` Integer to keep track of what category code table to allocate. In LuaTeX it is only used in format mode to implement `\cctab_new:N`. In other engines it is used to make csnames for dynamic tables.

```

22657 \int_new:N \g__cctab_allocate_int

```

(End definition for `\g__cctab_allocate_int`.)

`\l__cctab_internal_a_tl` Scratch space. For instance, when popping `\g__cctab_stack_seq`/`\g__cctab_unused_seq`, consists of the catcodetable number (integer denotation) in LuaTeX, or of an intarray variable (as a single token) in other engines.

```

22658 \tl_new:N \l__cctab_internal_a_tl

```

```

22659 \tl_new:N \l__cctab_internal_b_tl

```

(End definition for `\l__cctab_internal_a_tl` and `\l__cctab_internal_b_tl`.)

`\g__cctab_endlinechar_prop` In LuaTeX we store the `\endlinechar` associated to each `\catcodetable` in a property list, unless it is the default value 13.

```

22660 \prop_new:N \g__cctab_endlinechar_prop

```

(End definition for `\g__cctab_endlinechar_prop`.)

38.2 Allocating category code tables

`\cctab_new:N` The `__cctab_new:N` auxiliary allocates a new catcode table but does not attempt to set its value consistently across engines. It is used both in `\cctab_new:N`, which sets catcodes to iniTeX values, and in `\cctab_begin:N`/`\cctab_end:` for dynamically allocated tables.

`__cctab_new:c` First, the LuaTeX case. Creating a new category code table is done like other registers. In ConTeXt, `\newcatcodetable` does not include the initialisation, so that is added explicitly.

`__cctab_new:N`

`__cctab_gstore:Nnn`

```

22661 \sys_if_engine luatex:TF
22662 {
22663   \cs_new_protected:Npn \cctab_new:N #1
22664   {
22665     \__kernel_chk_if_free_cs:N #1
22666     \__cctab_new:N #1
22667   }
22668   \cs_new_protected:Npn \__cctab_new:N #1
22669   {
22670     \newcatcodetable #1
22671     \tex_initcatcodetable:D #1
22672   }
22673 }

```

Now the case for other engines. Here, each table is an integer array. Following the LuaTeX pattern, a new table starts with `initEX` codes. The index base is out-by-one, so we have an internal function to handle that. The `initEX \endlinechar` is 13.

```

22674 {
22675   \cs_new_protected:Npn \__cctab_new:N #1
22676     { \intarray_new:Nn #1 { 257 } }
22677   \cs_new_protected:Npn \__cctab_gstore:Nnn #1#2#3
22678     { \intarray_gset:Nnn #1 { \int_eval:n { #2 + 1 } } {#3} }
22679   \cs_new_protected:Npn \cctab_new:N #1
22680     {
22681       \__kernel_chk_if_free_cs:N #1
22682       \__cctab_new:N #1
22683       \int_step_inline:nn { 256 }
22684         { \__kernel_intarray_gset:Nnn #1 {##1} { 12 } }
22685       \__kernel_intarray_gset:Nnn #1 { 257 } { 13 }
22686       \__cctab_gstore:Nnn #1 { 0 } { 9 }
22687       \__cctab_gstore:Nnn #1 { 13 } { 5 }
22688       \__cctab_gstore:Nnn #1 { 32 } { 10 }
22689       \__cctab_gstore:Nnn #1 { 37 } { 14 }
22690       \int_step_inline:nnn { 65 } { 90 }
22691         { \__cctab_gstore:Nnn #1 {##1} { 11 } }
22692       \__cctab_gstore:Nnn #1 { 92 } { 0 }
22693       \int_step_inline:nnn { 97 } { 122 }
22694         { \__cctab_gstore:Nnn #1 {##1} { 11 } }
22695       \__cctab_gstore:Nnn #1 { 127 } { 15 }
22696     }
22697   }
22698   \cs_generate_variant:Nn \cctab_new:N { c }

```

(End definition for `\cctab_new:N`, `__cctab_new:N`, and `__cctab_gstore:Nnn`. This function is documented on page 225.)

38.3 Saving category code tables

`__cctab_gset:n`
`__cctab_gset_aux:n`

In various functions we need to save the current catcodes (globally) in a table. In LuaTeX, saving the catcodes is a primitives, but the `\endlinechar` needs more work: to avoid filling `\g__cctab_endlinechar_prop` with many entries we special-case the default value 13. In other engines we store 256 current catcodes and the `\endlinechar` in an intarray variable.

```

22699 \sys_if_engine luatex:TF
22700 {
22701   \cs_new_protected:Npn \__cctab_gset:n #1
22702     { \exp_args:Nf \__cctab_gset_aux:n { \int_eval:n {#1} } }
22703   \cs_new_protected:Npn \__cctab_gset_aux:n #1
22704     {
22705       \tex_savecatcodetable:D #1 \scan_stop:
22706       \int_compare:nNnTF { \tex_endlinechar:D } = { 13 }
22707         { \prop_gremove:Nn \g__cctab_endlinechar_prop {#1} }
22708         {
22709           \prop_gput:NnV \g__cctab_endlinechar_prop {#1}
22710             \tex_endlinechar:D
22711         }
22712     }

```

```

22713 }
22714 {
22715   \cs_new_protected:Npn \__cctab_gset:n #1
22716   {
22717     \int_step_inline:nn { 256 }
22718     {
22719       \__kernel_intarray_gset:Nnn #1 {##1}
22720       { \char_value_catcode:n { ##1 - 1 } }
22721     }
22722     \__kernel_intarray_gset:Nnn #1 { 257 }
22723     { \tex_endlinechar:D }
22724   }
22725 }

```

(End definition for `__cctab_gset:n` and `__cctab_gset_aux:n`.)

`\cctab_gset:Nn` Category code tables are always global, so only one version of assignments is needed.
`\cctab_gset:cn` Simply run the setup in a group and save the result in a category code table #1, provided it is valid. The internal function is defined above depending on the engine.

```

22726 \cs_new_protected:Npn \cctab_gset:Nn #1#2
22727 {
22728   \__cctab_chk_if_valid:NT #1
22729   {
22730     \group_begin:
22731     \cctab_select:N \c_initex_cctab
22732     #2 \scan_stop:
22733     \__cctab_gset:n {#1}
22734     \group_end:
22735   }
22736 }
22737 \cs_generate_variant:Nn \cctab_gset:Nn { c }

```

(End definition for `\cctab_gset:Nn`. This function is documented on page 225.)

38.4 Using category code tables

`\g__cctab_internal_cctab` In LuaTeX, we must ensure that the saved tables are read-only. This is done by applying the saved table, then switching immediately to a scratch table. Any later catcode assignment will affect that scratch table rather than the saved one. If we simply switched to the saved tables, then `\char_set_catcode_other:N` in the example below would change `\c_document_cctab` and a later use of that table would give the wrong category code to `_`.

```

\use:n
{
  \cctab_begin:N \c_document_cctab
  \char_set_catcode_other:N \_
  \cctab_end:
  \cctab_begin:N \c_document_cctab
  \int_compare:nTF { \char_value_catcode:n { '_' } = 8 }
  { \TRUE } { \ERROR }
  \cctab_end:
}

```


We must also make sure that a scratch table is never reused in a nested group: in the following example, the scratch table used by the first `\cctab_begin:N` would be changed globally by the second one issuing `\savecatcodetable`, and after `\group_end:` the wrong category codes (those of `\c_str_cctab`) would be imposed. Note that the inner `\cctab_end:` restores the correct catcodes only locally, so the problem really comes up because of the different grouping level. The simplest is to use a scratch table labeled by the `\currentgrouplevel`. We initialize one of them as an example.

```
\use:n
{
  \cctab_begin:N \c_document_cctab
  \group_begin:
    \cctab_begin:N \c_str_cctab
    \cctab_end:
  \group_end:
  \cctab_end:
}

22738 \sys_if_engine luatex:T
22739 {
22740   \__cctab_new:N \g__cctab_internal_cctab
22741   \cs_new:Npn \__cctab_internal_cctab_name:
22742     {
22743       g__cctab_internal
22744       \tex_romannumeral:D \tex_currentgrouplevel:D
22745       _cctab
22746     }
22747 }
```

(End definition for `\g__cctab_internal_cctab` and `__cctab_internal_cctab_name:.`)

`\cctab_select:N` The public function simply checks the `\cctab var` exists before using the engine-dependent `__cctab_select:N`. Skipping these checks would result in low-level engine-dependent errors. First, the LuaTeX case. In other engines, selecting a catcode table is a matter of doing 256 catcode assignments and setting the `\endlinechar`.

```
22748 \cs_new_protected:Npn \cctab_select:N #1
22749 { \__cctab_chk_if_valid:NT #1 { \__cctab_select:N #1 } }
22750 \cs_generate_variant:Nn \cctab_select:N { c }
22751 \sys_if_engine luatex:TF
22752 {
22753   \cs_new_protected:Npn \__cctab_select:N #1
22754   {
22755     \tex_catcodetable:D #1
22756     \prop_get:NVNTF \g__cctab_endlinechar_prop #1 \l__cctab_internal_a_tl
22757     { \int_set:Nn \tex_endlinechar:D { \l__cctab_internal_a_tl } }
22758     { \int_set:Nn \tex_endlinechar:D { 13 } }
22759     \cs_if_exist:cF { \__cctab_internal_cctab_name: }
22760     { \exp_args:Nc \__cctab_new:N { \__cctab_internal_cctab_name: } }
22761     \exp_args:Nc \tex_savecatcodetable:D { \__cctab_internal_cctab_name: }
22762     \exp_args:Nc \tex_catcodetable:D { \__cctab_internal_cctab_name: }
22763   }
22764 }
22765 {
```

```

22766 \cs_new_protected:Npn \__cctab_select:N #1
22767 {
22768   \int_step_inline:nn { 256 }
22769   {
22770     \char_set_catcode:nn { ##1 - 1 }
22771     { \__kernel_intarray_item:Nn #1 {##1} }
22772   }
22773   \int_set:Nn \tex_endlinechar:D
22774   { \__kernel_intarray_item:Nn #1 { 257 } }
22775 }
22776 }

```

(End definition for `\cctab_select:N` and `__cctab_select:N`. This function is documented on page 225.)

`\g__cctab_next_cctab`
`__cctab_begin_aux:`

For `\cctab_begin:N/\cctab_end:` we will need to allocate dynamic tables. This is done here by `__cctab_begin_aux:`, which puts a table number (in LuaTeX) or name (in other engines) into `\l__cctab_internal_a_tl`. In LuaTeX this simply calls `__cctab_new:N` and uses the resulting catcodetable number; in other engines we need to give a name to the intarray variable and use that. In LuaTeX, to restore catcodes at `\cctab_end:` we cannot just set `\catcodetable` to its value before `\cctab_begin:N`, because that table may have been altered by other code in the mean time. So we must make sure to save the catcodes in a table we control and restore them at `\cctab_end:`.

```

22777 \sys_if_engine luatex:TF
22778 {
22779   \cs_new_protected:Npn \__cctab_begin_aux:
22780   {
22781     \__cctab_new:N \g__cctab_next_cctab
22782     \tl_set:NV \l__cctab_internal_a_tl \g__cctab_next_cctab
22783     \cs_undefine:N \g__cctab_next_cctab
22784   }
22785 }
22786 {
22787   \cs_new_protected:Npn \__cctab_begin_aux:
22788   {
22789     \int_gincr:N \g__cctab_allocate_int
22790     \exp_args:Nc \__cctab_new:N
22791     { g__cctab_ \int_use:N \g__cctab_allocate_int _cctab }
22792     \exp_args:NNc \tl_set:Nn \l__cctab_internal_a_tl
22793     { g__cctab_ \int_use:N \g__cctab_allocate_int _cctab }
22794   }
22795 }

```

(End definition for `\g__cctab_next_cctab` and `__cctab_begin_aux:`.)

`\cctab_begin:N`
`\cctab_begin:c`

Check the `\cctab var` exists, to avoid low-level errors. Get in `\l__cctab_internal_a_tl` the number/name of a dynamic table, either from `\g__cctab_unused_seq` where we save tables that are not currently in use, or from `__cctab_begin_aux:` if none are available. Then save the current catcodes into the table (pointed to by) `\l__cctab_internal_a_tl` and save that table number in a stack before selecting the desired catcodes.

```

22796 \cs_new_protected:Npn \cctab_begin:N #1
22797 {
22798   \__cctab_chk_if_valid:NT #1
22799   {

```

```

22800     \seq_gpop:N NF \g__cctab_unused_seq \l__cctab_internal_a_tl
22801     { \__cctab_begin_aux: }
22802     \exp_args:Nx \__cctab_chk_group_begin:n
22803     { \__cctab_nesting_number:N \l__cctab_internal_a_tl }
22804     \seq_gpush:NV \g__cctab_stack_seq \l__cctab_internal_a_tl
22805     \exp_args:NV \__cctab_gset:n \l__cctab_internal_a_tl
22806     \__cctab_select:N #1
22807   }
22808 }
22809 \cs_generate_variant:Nn \cctab_begin:N { c }

```

(End definition for `\cctab_begin:N`. This function is documented on page 225.)

\cctab_end: Make sure a `\cctab_begin:N` was used some time earlier, get in `\l__cctab_internal_a_tl` the catcode table number/name in which the prevailing catcodes were stored, then restore these catcodes. The dynamic table is now unused hence stored in `\g__cctab_unused_seq` for recycling by later `\cctab_begin:N`.

```

22810 \cs_new_protected:Npn \cctab_end:
22811 {
22812   \seq_gpop:NNTF \g__cctab_stack_seq \l__cctab_internal_a_tl
22813   {
22814     \seq_gpush:NV \g__cctab_unused_seq \l__cctab_internal_a_tl
22815     \exp_args:Nx \__cctab_chk_group_end:n
22816     { \__cctab_nesting_number:N \l__cctab_internal_a_tl }
22817     \__cctab_select:N \l__cctab_internal_a_tl
22818   }
22819   { \__kernel_msg_error:nn { kernel } { cctab-extra-end } }
22820 }

```

(End definition for `\cctab_end:`. This function is documented on page 225.)

`__cctab_chk_group_begin:n` `__cctab_chk_group_end:n` Catcode tables are not allowed to be intermixed with groups, so here we check that they are properly nested regarding \TeX groups. `__cctab_chk_group_begin:n` stores the current group level in a stack, and locally defines a dummy control sequence `__cctab_group_⟨cctab-level⟩_chk:`.

`__cctab_chk_group_end:n` pops the stack, and compares the returned value with `\tex_currentgrouplevel:D`. If they differ, `\cctab_end:` is in a different grouping level than the matching `\cctab_begin:N`. If they are the same, both happened at the same level, however a group might have ended and another started between `\cctab_begin:N` and `\cctab_end:`:

```

\group_begin:
  \cctab_begin:N \c_document_cctab
\group_end:
\group_begin:
  \cctab_end:
\group_end:

```

In this case checking `\tex_currentgrouplevel:D` is not enough, so we locally define `__cctab_group_⟨cctab-level⟩_chk:`, and then check if it exist in `\cctab_end:`. If it doesn't, we know there was a group end where it shouldn't.

The `⟨cctab-level⟩` in the sentinel macro above cannot be replaced by the more convenient `\tex_currentgrouplevel:D` because with the latter we might be tricked. Suppose:

```

\group_begin:
  \cctab_begin:N \c_code_cctab % A
\group_end:
\group_begin:
  \cctab_begin:N \c_code_cctab % B
  \cctab_end: % C
  \cctab_end: % D
\group_end:

```

The line marked with A would start a `cctab` with a sentinel token named `__cctab_group_1_chk:`, which would disappear at the `\group_end:` that follows. But B would create the same sentinel token, since both are at the same group level. Line C would end the `cctab` from line B correctly, but so would line D because line B created the same sentinel token. Using `\cctab_level` works correctly because it signals that certain `cctab` level was activated somewhere, but if it doesn't exist when the `\cctab_end:` is reached, we had a problem.

Unfortunately these tests only flag the wrong usage at the `\cctab_end:`, which might be far from the `\cctab_begin:N`. However it isn't possible to signal the wrong usage at the `\group_end:` without using `\tex_aftergroup:D`, which is unsafe in certain types of groups.

The three cases checked here just raise an error, and no recovery is attempted: usually interleaving groups and catcode tables will work predictably.

```

22821 \cs_new_protected:Npn \__cctab_chk_group_begin:n #1
22822 {
22823   \seq_gpush:Nx \g__cctab_group_seq
22824   { \int_use:N \tex_currentgrouplevel:D }
22825   \cs_set_eq:cN { __cctab_group_ #1 _chk: } \prg_do_nothing:
22826 }
22827 \cs_new_protected:Npn \__cctab_chk_group_end:n #1
22828 {
22829   \seq_gpop:NN \g__cctab_group_seq \l__cctab_internal_b_tl
22830   \bool_lazy_and:nnF
22831   {
22832     \int_compare_p:nNn
22833     { \tex_currentgrouplevel:D } = { \l__cctab_internal_b_tl }
22834   }
22835   { \cs_if_exist_p:c { __cctab_group_ #1 _chk: } }
22836   {
22837     \__kernel_msg_error:nxx { kernel } { cctab-group-mismatch }
22838     {
22839       \int_sign:n
22840       { \tex_currentgrouplevel:D - \l__cctab_internal_b_tl }
22841     }
22842   }
22843   \cs_undefine:c { __cctab_group_ #1 _chk: }
22844 }

```

(End definition for `__cctab_chk_group_begin:n` and `__cctab_chk_group_end:n`.)

```

\__cctab_nesting_number:N
\__cctab_nesting_number:w

```

This macro returns the numeric index of the current catcode table. In LuaTeX this is just the argument, which is a count reference to a `\catcodetable` register. In other engines, the number is extracted from the `cctab` variable.

```

22845 \sys_if_engine luatex:TF
22846 { \cs_new:Npn \__cctab_nesting_number:N #1 {#1} }
22847 {
22848   \cs_new:Npn \__cctab_nesting_number:N #1
22849   {
22850     \exp_after:wN \exp_after:wN \exp_after:wN \__cctab_nesting_number:w
22851     \exp_after:wN \token_to_str:N #1
22852   }
22853   \use:x
22854   {
22855     \cs_new:Npn \exp_not:N \__cctab_nesting_number:w
22856     ##1 \tl_to_str:n { g__cctab_ } ##2 \tl_to_str:n { _cctab } {##2}
22857   }
22858 }

```

(End definition for `__cctab_nesting_number:N` and `__cctab_nesting_number:w`.)

Finally, install some code at the end of the \TeX run to check that all `\cctab_begin:N` were ended by some `\cctab_end:`.

```

22859 \cs_if_exist:NT \hook_gput_code:nnn
22860 {
22861   \hook_gput_code:nnn { enddocument/end } { kernel }
22862   {
22863     \seq_if_empty:NF \g__cctab_stack_seq
22864     { \__kernel_msg_error:nn { kernel } { cctab-missing-end } }
22865   }
22866 }

```

38.5 Category code table conditionals

`\cctab_if_exist:N` Checks whether a $\langle cctab\ var \rangle$ is defined.

```

\cctab_if_exist:c
22867 \prg_new_eq_conditional:NNn \cctab_if_exist:N \cs_if_exist:N
22868 { TF , T , F , p }
22869 \prg_new_eq_conditional:NNn \cctab_if_exist:c \cs_if_exist:c
22870 { TF , T , F , p }

```

(End definition for `\cctab_if_exist:N`. This function is documented on page ??.)

`__cctab_chk_if_valid:NTF` Checks whether the argument is defined and whether it is a valid $\langle cctab\ var \rangle$. In Lua \TeX the validity of the $\langle cctab\ var \rangle$ is checked by the engine, which complains if the argument is not a `\chardef`'ed constant. In other engines, check if the given command is an intarray variable (the underlying definition is a copy of the `cmr10` font).

```

22871 \prg_new_protected_conditional:Npnn \__cctab_chk_if_valid:N #1
22872 { TF , T , F }
22873 {
22874   \cctab_if_exist:NTF #1
22875   {
22876     \__cctab_chk_if_valid_aux:NTF #1
22877     { \prg_return_true: }
22878     {
22879       \__kernel_msg_error:nnx { kernel } { invalid-cctab }
22880       { \token_to_str:N #1 }
22881       \prg_return_false:
22882     }

```

```

22883     }
22884     {
22885         \_kernel_msg_error:nxx { kernel } { command-not-defined }
22886         { \token_to_str:N #1 }
22887         \prg_return_false:
22888     }
22889 }
22890 \sys_if_engine luatex:TF
22891 {
22892     \cs_new_protected:Npn \__cctab_chk_if_valid_aux:NTF #1
22893     {
22894         \int_compare:nNnTF {#1-1} < { \e@alloc@ccodetable@count }
22895     }
22896     \cs_if_exist:NT \c_syst_catcodes_n
22897     {
22898         \cs_gset_protected:Npn \__cctab_chk_if_valid_aux:NTF #1
22899         {
22900             \int_compare:nTF { #1 <= \c_syst_catcodes_n }
22901         }
22902     }
22903 }
22904 {
22905     \cs_new_protected:Npn \__cctab_chk_if_valid_aux:NTF #1
22906     {
22907         \exp_args:Nf \str_if_in:nnTF
22908         { \cs_meaning:N #1 }
22909         { select~font~cmr10~at~ }
22910     }
22911 }

```

(End definition for __cctab_chk_if_valid:NTF and __cctab_chk_if_valid_aux:NTF.)

38.6 Constant category code tables

\cctab_const:Nn Creates a new *cctab var* then sets it with the current and user-supplied codes.

```

\cctab_const:cn
22912 \cs_new_protected:Npn \cctab_const:Nn #1#2
22913 {
22914     \cctab_new:N #1
22915     \cctab_gset:Nn #1 {#2}
22916 }
22917 \cs_generate_variant:Nn \cctab_const:Nn { c }

```

(End definition for \cctab_const:Nn. This function is documented on page 225.)

\c_initex_cctab Creating category code tables means thinking starting from $\text{iniT}_{\text{E}}\text{X}$. For all-other and
\c_other_cctab the standard “string” tables that’s easy.

```

\c_str_cctab
22918 \cctab_new:N \c_initex_cctab
22919 \cctab_const:Nn \c_other_cctab
22920 {
22921     \cctab_select:N \c_initex_cctab
22922     \int_set:Nn \tex_endlinechar:D { -1 }
22923     \int_step_inline:nnn { 0 } { 127 }
22924     { \char_set_catcode_other:n {#1} }
22925 }

```

```

22926 \cctab_const:Nn \c_str_cctab
22927 {
22928   \cctab_select:N \c_other_cctab
22929   \char_set_catcode_space:n { 32 }
22930 }

```

(End definition for `\c_initex_cctab`, `\c_other_cctab`, and `\c_str_cctab`. These variables are documented on page 226.)

`\c_code_cctab`
`\c_document_cctab`

To pick up document-level category codes, we need to delay set up to the end of the format, where that's possible. Also, as there are a *lot* of category codes to set, we avoid using the official interface and store the document codes using internal code. Depending on whether we are in the hook or not, the catcodes may be code or document, so we explicitly set up both correctly.

```

22931 \cs_if_exist:NTF \@expl@finalise@setup@@
22932 { \tl_gput_right:Nn \@expl@finalise@setup@@ }
22933 { \use:n }
22934 {
22935   \__cctab_new:N \c_code_cctab
22936   \group_begin:
22937     \int_set:Nn \tex_endlinechar:D { 32 }
22938     \char_set_catcode_invalid:n { 0 }
22939     \bool_lazy_or:nnTF
22940       { \sys_if_engine_xetex_p: } { \sys_if_engine_luatex_p: }
22941       { \int_step_function:nnN { 31 } { 64 } \char_set_catcode_invalid:n }
22942       { \int_step_function:nnN { 31 } { 64 } \char_set_catcode_active:n }
22943     \int_step_function:nnN { 33 } { 64 } \char_set_catcode_other:n
22944     \int_step_function:nnN { 65 } { 90 } \char_set_catcode_letter:n
22945     \int_step_function:nnN { 91 } { 96 } \char_set_catcode_other:n
22946     \int_step_function:nnN { 97 } { 122 } \char_set_catcode_letter:n
22947     \char_set_catcode_ignore:n { 9 } % tab
22948     \char_set_catcode_other:n { 10 } % lf
22949     \char_set_catcode_active:n { 12 } % ff
22950     \char_set_catcode_end_line:n { 13 } % cr
22951     \char_set_catcode_ignore:n { 32 } % space
22952     \char_set_catcode_parameter:n { 35 } % hash
22953     \char_set_catcode_math_toggle:n { 36 } % dollar
22954     \char_set_catcode_comment:n { 37 } % percent
22955     \char_set_catcode_alignment:n { 38 } % ampersand
22956     \char_set_catcode_letter:n { 58 } % colon
22957     \char_set_catcode_escape:n { 92 } % backslash
22958     \char_set_catcode_math_superscript:n { 94 } % circumflex
22959     \char_set_catcode_letter:n { 95 } % underscore
22960     \char_set_catcode_group_begin:n { 123 } % left brace
22961     \char_set_catcode_other:n { 124 } % pipe
22962     \char_set_catcode_group_end:n { 125 } % right brace
22963     \char_set_catcode_space:n { 126 } % tilde
22964     \char_set_catcode_invalid:n { 127 } % ^^?
22965     \bool_lazy_or:nnF
22966       { \sys_if_engine_xetex_p: } { \sys_if_engine_luatex_p: }
22967       { \int_step_function:nnN { 128 } { 255 } \char_set_catcode_active:n }
22968     \__cctab_gset:n { \c_code_cctab }
22969   \group_end:
22970   \cctab_const:Nn \c_document_cctab

```

```

22971     {
22972         \cctab_select:N \c_code_cctab
22973         \int_set:Nn \tex_endlinechar:D { 13 }
22974         \char_set_catcode_space:n      { 9 }
22975         \char_set_catcode_space:n      { 32 }
22976         \char_set_catcode_other:n      { 58 }
22977         \char_set_catcode_math_subscript:n { 95 }
22978         \char_set_catcode_active:n     { 126 }
22979     }
22980 }

```

(End definition for `\c_code_cctab` and `\c_document_cctab`. These variables are documented on page 226.)

38.7 Messages

```

22981 \__kernel_msg_new:nnnn { kernel } { cctab-stack-full }
22982 { The-category-code-table-stack-is-exhausted. }
22983 {
22984     LaTeX-has-been-asked-to-switch-to-a-new-category-code-table,~
22985     but-there-is-no-more-space-to-do-this!
22986 }
22987 \__kernel_msg_new:nnnn { kernel } { cctab-extra-end }
22988 { Extra~\iow_char:N\\cctab_end:~ignored~\msg_line_context:. }
22989 {
22990     LaTeX-came-across-a~\iow_char:N\\cctab_end:~without-a-matching~
22991     \iow_char:N\\cctab_begin:N.~This-command-will-be-ignored.
22992 }
22993 \__kernel_msg_new:nnnn { kernel } { cctab-missing-end }
22994 { Missing~\iow_char:N\\cctab_end:~before-end-of-TeX-run. }
22995 {
22996     LaTeX-came-across-more~\iow_char:N\\cctab_begin:N~than~
22997     \iow_char:N\\cctab_end:.
22998 }
22999 \__kernel_msg_new:nnnn { kernel } { invalid-cctab }
23000 { Invalid~\iow_char:N\\catcode~table. }
23001 {
23002     You-can-only-switch-to-a~\iow_char:N\\catcode~table-that-is~
23003     initialized-using~\iow_char:N\\cctab_new:N~or~
23004     \iow_char:N\\cctab_const:Nn.
23005 }
23006 \__kernel_msg_new:nnnn { kernel } { cctab-group-mismatch }
23007 {
23008     \iow_char:N\\cctab_end:~occurred-in-a~
23009     \int_case:nn {#1}
23010     {
23011         { 0 } { different-group }
23012         { 1 } { higher-group-level }
23013         { -1 } { lower-group-level }
23014     } ~than~
23015     the-matching~\iow_char:N\\cctab_begin:N.
23016 }
23017 {
23018     Catcode-tables-and-groups-must-be-properly-nested,~but~

```



```

23019     you~tried~to~interleave~them.~LaTeX~will~try~to~proceed,~
23020     but~results~may~be~unexpected.
23021   }
23022 </package>

```

39 l3sort implementation

```

23023 <*package>
23024 <@@=sort>

```

39.1 Variables

`\g__sort_internal_seq` Sorting happens in a group; the result is stored in those global variables before being copied outside the group to the proper places. For `seq` and `tl` this is more efficient than using `\use:x` (or some `\exp_args:NNNx`) to smuggle the definition outside the group since TeX does not need to re-read tokens. For `clist` we don't gain anything since the result is converted from `seq` to `clist` anyways.

```

23025 \seq_new:N \g__sort_internal_seq
23026 \tl_new:N \g__sort_internal_tl

```

(End definition for `\g__sort_internal_seq` and `\g__sort_internal_tl`.)

`\l__sort_length_int` The sequence has `\l__sort_length_int` items and is stored from `\l__sort_min_int` to `\l__sort_top_int - 1`. While reading the sequence in memory, we check that `\l__sort_top_int` remains at most `\l__sort_max_int`, precomputed by `__sort_compute_range:.` That bound is such that the merge sort only uses `\toks` registers less than `\l__sort_true_max_int`, namely those that have not been allocated for use in other code: the user's comparison code could alter these.

```

23027 \int_new:N \l__sort_length_int
23028 \int_new:N \l__sort_min_int
23029 \int_new:N \l__sort_top_int
23030 \int_new:N \l__sort_max_int
23031 \int_new:N \l__sort_true_max_int

```

(End definition for `\l__sort_length_int` and others.)

`\l__sort_block_int` Merge sort is done in several passes. In each pass, blocks of size `\l__sort_block_int` are merged in pairs. The block size starts at 1, and, for a length in the range $[2^k + 1, 2^{k+1}]$, reaches 2^k in the last pass.

```

23032 \int_new:N \l__sort_block_int

```

(End definition for `\l__sort_block_int`.)

`\l__sort_begin_int` When merging two blocks, `\l__sort_begin_int` marks the lowest index in the two blocks, and `\l__sort_end_int` marks the highest index, plus 1.

```

23033 \int_new:N \l__sort_begin_int
23034 \int_new:N \l__sort_end_int

```

(End definition for `\l__sort_begin_int` and `\l__sort_end_int`.)

`\l__sort_A_int` When merging two blocks (whose end-points are `beg` and `end`), A starts from the high end of the low block, and decreases until reaching `beg`. The index B starts from the top of the range and marks the register in which a sorted item should be put. Finally, C points to the copy of the high block in the interval of registers starting at `\l__sort_length_int`, upwards. C starts from the upper limit of that range.

```
23035 \int_new:N \l__sort_A_int
23036 \int_new:N \l__sort_B_int
23037 \int_new:N \l__sort_C_int
```

(End definition for `\l__sort_A_int`, `\l__sort_B_int`, and `\l__sort_C_int`.)

`\s__sort_mark` Internal scan marks.

```
\s__sort_stop 23038 \scan_new:N \s__sort_mark
23039 \scan_new:N \s__sort_stop
```

(End definition for `\s__sort_mark` and `\s__sort_stop`.)

39.2 Finding available `\toks` registers

`__sort_shrink_range:` After `__sort_compute_range:` (defined below) determines that `\toks` registers between `\l__sort_min_int` (included) and `\l__sort_true_max_int` (excluded) have not yet been assigned, `__sort_shrink_range:` computes `\l__sort_max_int` to reflect the need for a buffer when merging blocks in the merge sort. Given $2^n \leq A \leq 2^n + 2^{n-1}$ registers we can sort $\lfloor A/2 \rfloor + 2^{n-2}$ items while if we have $2^n + 2^{n-1} \leq A \leq 2^{n+1}$ registers we can sort $A - 2^{n-1}$ items. We first find out a power 2^n such that $2^n \leq A \leq 2^{n+1}$ by repeatedly halving `\l__sort_block_int`, starting at 2^{15} or 2^{14} namely half the total number of registers, then we use the formulas and set `\l__sort_max_int`.

```
23040 \cs_new_protected:Npn \__sort_shrink_range:
23041 {
23042   \int_set:Nn \l__sort_A_int
23043     { \l__sort_true_max_int - \l__sort_min_int + 1 }
23044   \int_set:Nn \l__sort_block_int { \c_max_register_int / 2 }
23045   \__sort_shrink_range_loop:
23046   \int_set:Nn \l__sort_max_int
23047     {
23048       \int_compare:nNnTF
23049         { \l__sort_block_int * 3 / 2 } > \l__sort_A_int
23050         {
23051           \l__sort_min_int
23052           + ( \l__sort_A_int - 1 ) / 2
23053           + \l__sort_block_int / 4
23054           - 1
23055         }
23056         { \l__sort_true_max_int - \l__sort_block_int / 2 }
23057     }
23058 }
23059 \cs_new_protected:Npn \__sort_shrink_range_loop:
23060 {
23061   \if_int_compare:w \l__sort_A_int < \l__sort_block_int
23062     \tex_divide:D \l__sort_block_int 2 \exp_stop_f:
23063     \exp_after:wN \__sort_shrink_range_loop:
23064   \fi:
23065 }
```

(End definition for `__sort_shrink_range:` and `__sort_shrink_range_loop:.`)

`__sort_compute_range:` First find out what `\toks` have not yet been assigned. There are many cases. In $\text{\LaTeX} 2_{\epsilon}$ with no package, available `\toks` range from `\count15 + 1` to `\c_max_register_int` included (this was not altered despite the 2015 changes). When `\loctoks` is defined, namely in plain (e) \TeX , or when the package `etex` is loaded in $\text{\LaTeX} 2_{\epsilon}$, redefine `__sort_compute_range:` to use the range `\count265` to `\count275 - 1`. The `elocalloc` package also defines `\loctoks` but uses yet another number for the upper bound, namely `\e@alloc@top` (minus one). We must check for `\loctoks` every time a sorting function is called, as `etex` or `elocalloc` could be loaded.

In Con \TeX t MkIV the range is from `\c_syst_last_allocated_toks+1` to `\c_max_register_int`, and in MkII it is from `\lastallocatedtoks+1` to `\c_max_register_int`. In all these cases, call `__sort_shrink_range:.`

```

23066 \cs_new_protected:Npn \__sort_compute_range:
23067 {
23068   \int_set:Nn \l__sort_min_int { \tex_count:D 15 + 1 }
23069   \int_set:Nn \l__sort_true_max_int { \c_max_register_int + 1 }
23070   \__sort_shrink_range:
23071   \if_meaning:w \loctoks \tex_undefined:D \else:
23072     \if_meaning:w \loctoks \scan_stop: \else:
23073       \__sort_redefine_compute_range:
23074       \__sort_compute_range:
23075     \fi:
23076   \fi:
23077 }
23078 \cs_new_protected:Npn \__sort_redefine_compute_range:
23079 {
23080   \cs_if_exist:cTF { ver@elocalloc.sty }
23081   {
23082     \cs_gset_protected:Npn \__sort_compute_range:
23083     {
23084       \int_set:Nn \l__sort_min_int { \tex_count:D 265 }
23085       \int_set_eq:NN \l__sort_true_max_int \e@alloc@top
23086       \__sort_shrink_range:
23087     }
23088   }
23089   {
23090     \cs_gset_protected:Npn \__sort_compute_range:
23091     {
23092       \int_set:Nn \l__sort_min_int { \tex_count:D 265 }
23093       \int_set:Nn \l__sort_true_max_int { \tex_count:D 275 }
23094       \__sort_shrink_range:
23095     }
23096   }
23097 }
23098 \cs_if_exist:NT \loctoks { \__sort_redefine_compute_range: }
23099 \tl_map_inline:nn { \lastallocatedtoks \c_syst_last_allocated_toks }
23100 {
23101   \cs_if_exist:NT #1
23102   {
23103     \cs_gset_protected:Npn \__sort_compute_range:
23104     {
23105       \int_set:Nn \l__sort_min_int { #1 + 1 }

```

```

23106         \int_set:Nn \l__sort_true_max_int { \c_max_register_int + 1 }
23107         \__sort_shrink_range:
23108     }
23109 }
23110 }

```

(End definition for `__sort_compute_range:`, `__sort_redefine_compute_range:`, and `\c_sort_max_length_int`.)

39.3 Protected user commands

`__sort_main:NNNn` Sorting happens in three steps. First store items in `\toks` registers ranging from `\l__sort_min_int` to `\l__sort_top_int - 1`, while checking that the list is not too long. If we reach the maximum length, that's an error; exit the group. Secondly, sort the array of `\toks` registers, using the user-defined sorting function: `__sort_level:` calls `__sort_compare:nn` as needed. Finally, unpack the `\toks` registers (now sorted) into the target `tl`, or into `\g__sort_internal_seq` for `seq` and `clist`. This is done by `__sort_seq:NNNNn` and `__sort_tl:NNn`.

```

23111 \cs_new_protected:Npn \__sort_main:NNNn #1#2#3#4
23112 {
23113     \__sort_disable_toksdef:
23114     \__sort_compute_range:
23115     \int_set_eq:NN \l__sort_top_int \l__sort_min_int
23116     #1 #3
23117     {
23118         \if_int_compare:w \l__sort_top_int = \l__sort_max_int
23119             \__sort_too_long_error:NNw #2 #3
23120         \fi:
23121         \tex_toks:D \l__sort_top_int {##1}
23122         \int_incr:N \l__sort_top_int
23123     }
23124     \int_set:Nn \l__sort_length_int
23125     { \l__sort_top_int - \l__sort_min_int }
23126     \cs_set:Npn \__sort_compare:nn ##1 ##2 {#4}
23127     \int_set:Nn \l__sort_block_int { 1 }
23128     \__sort_level:
23129 }

```

(End definition for `__sort_main:NNNn`.)

`\tl_sort:Nn` Call the main sorting function then unpack `\toks` registers outside the group into the target token list. The unpacking is done by `__sort_tl_toks:w`; registers are numbered
`\tl_sort:cn` from `\l__sort_min_int` to `\l__sort_top_int - 1`. For expansion behaviour we need
`\tl_gsort:Nn` a couple of primitives. The `\tl_gclear:N` reduces memory usage. The `\prg_break_point:` is used by `__sort_main:NNNn` when the list is too long.
`\tl_gsort:cn`
`__sort_tl:NNn`
`__sort_tl_toks:w`

```

23130 \cs_new_protected:Npn \tl_sort:Nn { \__sort_tl:NNn \tl_set_eq:NN }
23131 \cs_generate_variant:Nn \tl_sort:Nn { c }
23132 \cs_new_protected:Npn \tl_gsort:Nn { \__sort_tl:NNn \tl_gset_eq:NN }
23133 \cs_generate_variant:Nn \tl_gsort:Nn { c }
23134 \cs_new_protected:Npn \__sort_tl:NNn #1#2#3
23135 {
23136     \group_begin:
23137     \__sort_main:NNNn \tl_map_inline:Nn \tl_map_break:n #2 {#3}

```

```

23138     \__kernel_tl_gset:Nx \g__sort_internal_tl
23139     { \__sort_tl_toks:w \l__sort_min_int ; }
23140   \group_end:
23141   #1 #2 \g__sort_internal_tl
23142   \tl_gclear:N \g__sort_internal_tl
23143   \prg_break_point:
23144 }
23145 \cs_new:Npn \__sort_tl_toks:w #1 ;
23146 {
23147   \if_int_compare:w #1 < \l__sort_top_int
23148   { \tex_the:D \tex_toks:D #1 }
23149   \exp_after:wN \__sort_tl_toks:w
23150   \int_value:w \int_eval:n { #1 + 1 } \exp_after:wN ;
23151   \fi:
23152 }

```

(End definition for `\tl_sort:Nn` and others. These functions are documented on page 53.)

```

\seq_sort:Nn Use the same general framework for seq and clist. Apply the general sorting code, then
\seq_sort:cn unpack \toks into \g__sort_internal_seq. Outside the group copy or convert (for
\seq_gsort:Nn clist) the data to the target variable. The \seq_gclear:N reduces memory usage. The
\seq_gsort:cn \prg_break_point: is used by \__sort_main:NNNn when the list is too long.
\clist_sort:Nn
\clist_sort:cn
\clist_gsort:Nn
\clist_gsort:cn
\__sort_seq:NNNNn
23153 \cs_new_protected:Npn \seq_sort:Nn
23154 { \__sort_seq:NNNNn \seq_map_inline:Nn \seq_map_break:n \seq_set_eq:NN }
23155 \cs_generate_variant:Nn \seq_sort:Nn { c }
23156 \cs_new_protected:Npn \seq_gsort:Nn
23157 { \__sort_seq:NNNNn \seq_map_inline:Nn \seq_map_break:n \seq_gset_eq:NN }
23158 \cs_generate_variant:Nn \seq_gsort:Nn { c }
23159 \cs_new_protected:Npn \clist_sort:Nn
23160 {
23161   \__sort_seq:NNNNn \clist_map_inline:Nn \clist_map_break:n
23162   \clist_set_from_seq:NN
23163 }
23164 \cs_generate_variant:Nn \clist_sort:Nn { c }
23165 \cs_new_protected:Npn \clist_gsort:Nn
23166 {
23167   \__sort_seq:NNNNn \clist_map_inline:Nn \clist_map_break:n
23168   \clist_gset_from_seq:NN
23169 }
23170 \cs_generate_variant:Nn \clist_gsort:Nn { c }
23171 \cs_new_protected:Npn \__sort_seq:NNNNn #1#2#3#4#5
23172 {
23173   \group_begin:
23174   \__sort_main:NNNn #1 #2 #4 {#5}
23175   \seq_gset_from_inline_x:Nnn \g__sort_internal_seq
23176   {
23177     \int_step_function:nnN
23178     { \l__sort_min_int } { \l__sort_top_int - 1 }
23179   }
23180   { \tex_the:D \tex_toks:D ##1 }
23181   \group_end:
23182   #3 #4 \g__sort_internal_seq
23183   \seq_gclear:N \g__sort_internal_seq
23184   \prg_break_point:

```

```
23185 }
```

(End definition for `\seq_sort:Nn` and others. These functions are documented on page 79.)

39.4 Merge sort

`__sort_level:` This function is called once blocks of size `\l__sort_block_int` (initially 1) are each sorted. If the whole list fits in one block, then we are done (this also takes care of the case of an empty list or a list with one item). Otherwise, go through pairs of blocks starting from 0, then double the block size, and repeat.

```
23186 \cs_new_protected:Npn \__sort_level:
23187 {
23188   \if_int_compare:w \l__sort_block_int < \l__sort_length_int
23189     \l__sort_end_int \l__sort_min_int
23190     \__sort_merge_blocks:
23191     \tex_advance:D \l__sort_block_int \l__sort_block_int
23192     \exp_after:wN \__sort_level:
23193   \fi:
23194 }
```

(End definition for `__sort_level:.`)

`__sort_merge_blocks:` This function is called to merge a pair of blocks, starting at the last value of `\l__sort_end_int` (end-point of the previous pair of blocks). If shifting by one block to the right we reach the end of the list, then this pass has ended: the end of the list is sorted already. Otherwise, store the result of that shift in *A*, which indexes the first block starting from the top end. Then locate the end-point (maximum) of the second block: shift *end* upwards by one more block, but keeping it \leq *top*. Copy this upper block of `\toks` registers in registers above *length*, indexed by *C*: this is covered by `__sort_copy_block:.` Once this is done we are ready to do the actual merger using `__sort_merge_blocks_aux:`, after shifting *A*, *B* and *C* so that they point to the largest index in their respective ranges rather than pointing just beyond those ranges. Of course, once that pair of blocks is merged, move on to the next pair.

```
23195 \cs_new_protected:Npn \__sort_merge_blocks:
23196 {
23197   \l__sort_begin_int \l__sort_end_int
23198   \tex_advance:D \l__sort_end_int \l__sort_block_int
23199   \if_int_compare:w \l__sort_end_int < \l__sort_top_int
23200     \l__sort_A_int \l__sort_end_int
23201     \tex_advance:D \l__sort_end_int \l__sort_block_int
23202     \if_int_compare:w \l__sort_end_int > \l__sort_top_int
23203       \l__sort_end_int \l__sort_top_int
23204     \fi:
23205     \l__sort_B_int \l__sort_A_int
23206     \l__sort_C_int \l__sort_top_int
23207     \__sort_copy_block:
23208     \int_decr:N \l__sort_A_int
23209     \int_decr:N \l__sort_B_int
23210     \int_decr:N \l__sort_C_int
23211     \exp_after:wN \__sort_merge_blocks_aux:
23212     \exp_after:wN \__sort_merge_blocks:
23213   \fi:
23214 }
```

(End definition for _sort_merge_blocks:.)

_sort_copy_block: We wish to store a copy of the “upper” block of \toks registers, ranging between the initial value of \l_sort_B_int (included) and \l_sort_end_int (excluded) into a new range starting at the initial value of \l_sort_C_int, namely \l_sort_top_int.

```

23215 \cs_new_protected:Npn \_sort_copy_block:
23216 {
23217   \tex_toks:D \l\_sort_C_int \tex_toks:D \l\_sort_B_int
23218   \int_incr:N \l\_sort_C_int
23219   \int_incr:N \l\_sort_B_int
23220   \if_int_compare:w \l\_sort_B_int = \l\_sort_end_int
23221     \use_i:nn
23222   \fi:
23223   \_sort_copy_block:
23224 }

```

(End definition for _sort_copy_block:.)

_sort_merge_blocks_aux: At this stage, the first block starts at \l_sort_begin_int, and ends at \l_sort_A_int, and the second block starts at \l_sort_top_int and ends at \l_sort_C_int. The result of the merger is stored at positions indexed by \l_sort_B_int, which starts at \l_sort_end_int – 1 and decreases down to \l_sort_begin_int, covering the full range of the two blocks. In other words, we are building the merger starting with the largest values. The comparison function is defined to return either **swapped** or **same**. Of course, this means the arguments need to be given in the order they appear originally in the list.

```

23225 \cs_new_protected:Npn \_sort_merge_blocks_aux:
23226 {
23227   \exp_after:wN \_sort_compare:nn \exp_after:wN
23228   { \tex_the:D \tex_toks:D \exp_after:wN \l\_sort_A_int \exp_after:wN }
23229   \exp_after:wN { \tex_the:D \tex_toks:D \l\_sort_C_int }
23230   \prg_do_nothing:
23231   \_sort_return_mark:w
23232   \_sort_return_mark:w
23233   \s\_sort_mark
23234   \_sort_return_none_error:
23235 }

```

(End definition for _sort_merge_blocks_aux:.)

\sort_return_same: Each comparison should call \sort_return_same: or \sort_return_swapped: exactly once. If neither is called, _sort_return_none_error: is called, since the **return_mark** removes tokens until \s_sort_mark. If one is called, the **return_mark** auxiliary removes everything except _sort_return_same:w (or its **swapped** analogue) followed by _sort_return_none_error:. Finally if two or more are called, _sort_return_two_error: ends up before any _sort_return_mark:w, so that it produces an error.

```

23236 \cs_new_protected:Npn \sort_return_same:
23237 #1 \_sort_return_mark:w #2 \s\_sort_mark
23238 {
23239   #1
23240   #2
23241   \_sort_return_two_error:
23242   \_sort_return_mark:w

```

```

23243     \s__sort_mark
23244     \__sort_return_same:w
23245   }
23246   \cs_new_protected:Npn \sort_return_swapped:
23247     #1 \__sort_return_mark:w #2 \s__sort_mark
23248   {
23249     #1
23250     #2
23251     \__sort_return_two_error:
23252     \__sort_return_mark:w
23253     \s__sort_mark
23254     \__sort_return_swapped:w
23255   }
23256   \cs_new_protected:Npn \__sort_return_mark:w #1 \s__sort_mark { }
23257   \cs_new_protected:Npn \__sort_return_none_error:
23258   {
23259     \__kernel_msg_error:nnxx { kernel } { return-none }
23260     { \tex_the:D \tex_toks:D \l__sort_A_int }
23261     { \tex_the:D \tex_toks:D \l__sort_C_int }
23262     \__sort_return_same:w \__sort_return_none_error:
23263   }
23264   \cs_new_protected:Npn \__sort_return_two_error:
23265   {
23266     \__kernel_msg_error:nnxx { kernel } { return-two }
23267     { \tex_the:D \tex_toks:D \l__sort_A_int }
23268     { \tex_the:D \tex_toks:D \l__sort_C_int }
23269   }

```

(End definition for `\sort_return_same:` and others. These functions are documented on page 227.)

`__sort_return_same:w` If the comparison function returns **same**, then the second argument fed to `__sort_compare:nn` should remain to the right of the other one. Since we build the merger starting from the right, we copy that `\toks` register into the allotted range, then shift the pointers *B* and *C*, and go on to do one more step in the merger, unless the second block has been exhausted: then the remainder of the first block is already in the correct registers and we are done with merging those two blocks.

```

23270   \cs_new_protected:Npn \__sort_return_same:w #1 \__sort_return_none_error:
23271   {
23272     \tex_toks:D \l__sort_B_int \tex_toks:D \l__sort_C_int
23273     \int_decr:N \l__sort_B_int
23274     \int_decr:N \l__sort_C_int
23275     \if_int_compare:w \l__sort_C_int < \l__sort_top_int
23276       \use_i:nn
23277     \fi:
23278     \__sort_merge_blocks_aux:
23279   }

```

(End definition for `__sort_return_same:w`.)

`__sort_return_swapped:w` If the comparison function returns **swapped**, then the next item to add to the merger is the first argument, contents of the `\toks` register *A*. Then shift the pointers *A* and *B* to the left, and go for one more step for the merger, unless the left block was exhausted (*A* goes below the threshold). In that case, all remaining `\toks` registers in the second block, indexed by *C*, are copied to the merger by `__sort_merge_blocks_end:`.


```

23280 \cs_new_protected:Npn \__sort_return_swapped:w #1 \__sort_return_none_error:
23281 {
23282   \tex_toks:D \l__sort_B_int \tex_toks:D \l__sort_A_int
23283   \int_decr:N \l__sort_B_int
23284   \int_decr:N \l__sort_A_int
23285   \if_int_compare:w \l__sort_A_int < \l__sort_begin_int
23286     \__sort_merge_blocks_end: \use_i:nn
23287   \fi:
23288   \__sort_merge_blocks_aux:
23289 }

```

(End definition for __sort_return_swapped:w.)

__sort_merge_blocks_end: This function’s task is to copy the \toks registers in the block indexed by C to the merger indexed by B . The end can equally be detected by checking when B reaches the threshold `begin`, or when C reaches `top`.

```

23290 \cs_new_protected:Npn \__sort_merge_blocks_end:
23291 {
23292   \tex_toks:D \l__sort_B_int \tex_toks:D \l__sort_C_int
23293   \int_decr:N \l__sort_B_int
23294   \int_decr:N \l__sort_C_int
23295   \if_int_compare:w \l__sort_B_int < \l__sort_begin_int
23296     \use_i:nn
23297   \fi:
23298   \__sort_merge_blocks_end:
23299 }

```

(End definition for __sort_merge_blocks_end:.)

39.5 Expandable sorting

Sorting expandably is very different from sorting and assigning to a variable. Since tokens cannot be stored, they must remain in the input stream, and be read through at every step. It is thus necessarily much slower (at best $O(n^2 \ln n)$) than non-expandable sorting functions ($O(n \ln n)$).

A prototypical version of expandable quicksort is as follows. If the argument has no item, return nothing, otherwise partition, using the first item as a pivot (argument `#4` of __sort:nnNnn). The arguments of __sort:nnNnn are 1. items less than `#4`, 2. items greater or equal to `#4`, 3. comparison, 4. pivot, 5. next item to test. If `#5` is the tail of the list, call \tl_sort:nN on `#1` and on `#2`, placing `#4` in between; \use:ff expands the parts to make \tl_sort:nN f-expandable. Otherwise, compare `#4` and `#5` using `#3`. If they are ordered, place `#5` amongst the “greater” items, otherwise amongst the “lesser” items, and continue partitioning.

```

\cs_new:Npn \tl_sort:nN #1#2
{
  \tl_if_blank:nF {#1}
  {
    \__sort:nnNnn { } { } #2
    #1 \q__sort_recursion_tail \q__sort_recursion_stop
  }
}
\cs_new:Npn \__sort:nnNnn #1#2#3#4#5

```

```

{
  \quark_if_recursion_tail_stop_do:nn {#5}
  { \use:ff { \tl_sort:nN {#1} #3 {#4} } { \tl_sort:nN {#2} #3 } }
  #3 {#4} {#5}
  { \__sort:nnNnn {#1} { #2 {#5} } #3 {#4} }
  { \__sort:nnNnn { #1 {#5} } {#2} #3 {#4} }
}
\cs_generate_variant:Nn \use:nn { ff }

```

There are quite a few optimizations available here: the code below is less legible, but more than twice as fast.

In the simple version of the code, `__sort:nnNnn` is called $O(n \ln n)$ times on average (the number of comparisons required by the quicksort algorithm). Hence most of our focus is on optimizing that function.

The first speed up is to avoid testing for the end of the list at every call to `__sort:nnNnn`. For this, the list is prepared by changing each *⟨item⟩* of the original token list into *⟨command⟩ {⟨item⟩}*, just like sequences are stored. We arrange things such that the *⟨command⟩* is the *⟨conditional⟩* provided by the user: the loop over the *⟨prepared tokens⟩* then looks like

```

\cs_new:Npn \__sort_loop:wNn ... #6#7
{
  #6 {⟨pivot⟩} {#7} ⟨loop big⟩ ⟨loop small⟩
  ⟨extra arguments⟩
}
\__sort_loop:wNn ... ⟨prepared tokens⟩
⟨end-loop⟩ {} \s__sort_stop

```

In this example, which matches the structure of `__sort_quick_split_i:NnnnnNn` and a few other functions below, the `__sort_loop:wNn` auxiliary normally receives the user's *⟨conditional⟩* as #6 and an *⟨item⟩* as #7. This is compared to the *⟨pivot⟩* (the argument #5, not shown here), and the *⟨conditional⟩* leaves the *⟨loop big⟩* or *⟨loop small⟩* auxiliary, which both have the same form as `__sort_loop:wNn`, receiving the next pair *⟨conditional⟩ {⟨item⟩}* as #6 and #7. At the end, #6 is the *⟨end-loop⟩* function, which terminates the loop.

The second speed up is to minimize the duplicated tokens between the **true** and **false** branches of the conditional. For this, we introduce two versions of `__sort:nnNnn`, which receive the new item as #1 and place it either into the list #2 of items less than the pivot #4 or into the list #3 of items greater or equal to the pivot.

```

\cs_new:Npn \__sort_i:nnnnNn #1#2#3#4#5#6
{
  #5 {#4} {#6} \__sort_ii:nnnnNn \__sort_i:nnnnNn
  {#6} { #2 {#1} } {#3} {#4}
}
\cs_new:Npn \__sort_ii:nnnnNn #1#2#3#4#5#6
{
  #5 {#4} {#6} \__sort_ii:nnnnNn \__sort_i:nnnnNn
  {#6} {#2} { #3 {#1} } {#4}
}

```

Note that the two functions have the form of `__sort_loop:wNn` above, receiving as `#5` the conditional or a function to end the loop. In fact, the lists `#2` and `#3` must be made of pairs $\langle conditional \rangle \{ \langle item \rangle \}$, so we have to replace `{#6}` above by `{ #5 {#6} }`, and `{#1}` by `#1`. The actual functions have one more argument, so all argument numbers are shifted compared to this code.

The third speed up is to avoid `\use:ff` using a continuation-passing style: `__sort_quick_split:NnNn` expects a list followed by `\s__sort_mark { \langle code \rangle }`, and expands to $\langle code \rangle \langle sorted list \rangle$. Sorting the two parts of the list around the pivot is done with

```
\__sort_quick_split:NnNn #2 ... \s__sort_mark
{
\__sort_quick_split:NnNn #1 ... \s__sort_mark { \langle code \rangle }
{ \langle pivot \rangle }
}
```

Items which are larger than the $\langle pivot \rangle$ are sorted, then placed after code that sorts the smaller items, and after the (braced) $\langle pivot \rangle$.

The fourth speed up is avoid the recursive call to `\tl_sort:nN` with an empty first argument. For this, we introduce functions similar to the `__sort_i:nnnnNn` of the last example, but aware of whether the list of $\langle conditional \rangle \{ \langle item \rangle \}$ read so far that are less than the pivot, and the list of those greater or equal, are empty or not: see `__sort_quick_split:NnNn` and functions defined below. Knowing whether the lists are empty or not is useless if we do not use distinct ending codes as appropriate. The splitting auxiliaries communicate to the $\langle end-loop \rangle$ function (that is initially placed after the “prepared” list) by placing a specific ending function, ignored when looping, but useful at the end. In fact, the $\langle end-loop \rangle$ function does nothing but place the appropriate ending function in front of all its arguments. The ending functions take care of sorting non-empty sublists, placing the pivot in between, and the continuation before.

The final change in fact slows down the code a little, but is required to avoid memory issues: schematically, when `TEX` encounters

```
\use:n { \use:n { \use:n { ... } ... } ... }
```

the argument of the first `\use:n` is not completely read by the second `\use:n`, hence must remain in memory; then the argument of the second `\use:n` is not completely read when grabbing the argument of the third `\use:n`, hence must remain in memory, and so on. The memory consumption grows quadratically with the number of nested `\use:n`. In practice, this means that we must read everything until a trailing `\s__sort_stop` once in a while, otherwise sorting lists of more than a few thousand items would exhaust a typical `TEX`’s memory.

`\tl_sort:nN`

`__sort_quick_prepare:Nnnn`

`__sort_quick_prepare_end:NNNnw`

`__sort_quick_cleanup:w`

The code within the `\exp_not:f` sorts the list, leaving in most cases a leading `\exp_not:f`, which stops the expansion, letting the result be return within `\exp_not:n`. We filter out the case of a list with no item, which would otherwise cause problems. Then prepare the token list `#1` by inserting the conditional `#2` before each item. The `prepare` auxiliary receives the conditional as `#1`, the prepared token list so far as `#2`, the next prepared item as `#3`, and the item after that as `#4`. The loop ends when `#4` contains `\prg_break_point:`, then the `prepare_end` auxiliary finds the prepared token list as `#4`. The scene is then set up for `__sort_quick_split:NnNn`, which sorts the prepared list

and perform the post action placed after `\s__sort_mark`, namely removing the trailing `\s__sort_stop` and `\s__sort_stop` and leaving `\exp_stop_f:` to stop f-expansion.

```

23300 \cs_new:Npn \tl_sort:nN #1#2
23301 {
23302   \exp_not:f
23303   {
23304     \tl_if_blank:nF {#1}
23305     {
23306       \__sort_quick_prepare:Nnnn #2 { } { }
23307       #1
23308       { \prg_break_point: \__sort_quick_prepare_end:NNNnw }
23309       \s__sort_stop
23310     }
23311   }
23312 }
23313 \cs_new:Npn \__sort_quick_prepare:Nnnn #1#2#3#4
23314 {
23315   \prg_break: #4 \prg_break_point:
23316   \__sort_quick_prepare:Nnnn #1 { #2 #3 } { #1 {#4} }
23317 }
23318 \cs_new:Npn \__sort_quick_prepare_end:NNNnw #1#2#3#4#5 \s__sort_stop
23319 {
23320   \__sort_quick_split:NnNn #4 \__sort_quick_end:nnTFNn { }
23321   \s__sort_mark { \__sort_quick_cleanup:w \exp_stop_f: }
23322   \s__sort_mark \s__sort_stop
23323 }
23324 \cs_new:Npn \__sort_quick_cleanup:w #1 \s__sort_mark \s__sort_stop {#1}

```

(End definition for `\tl_sort:nN` and others. This function is documented on page 53.)

`__sort_quick_split:NnNn` The `only_i`, `only_ii`, `split_i` and `split_ii` auxiliaries receive a useless first argument, the new item #2 (that they append to either one of the next two arguments), the list #3 of items less than the pivot, bigger items #4, the pivot #5, a *function* #6, and an item #7. The *function* is the user's *conditional* except at the end of the list where it is `__sort_quick_end:nnTFNn`. The comparison is applied to the *pivot* and the *item*, and calls the `only_i` or `split_i` auxiliaries if the *item* is smaller, and the `only_ii` or `split_ii` auxiliaries otherwise. In both cases, the next auxiliary goes to work right away, with no intermediate expansion that would slow down operations. Note that the argument #2 left for the next call has the form *conditional* {*item*}, so that the lists #3 and #4 keep the right form to be fed to the next sorting function. The `split` auxiliary differs from these in that it is missing three of the arguments, which would be empty, and its first argument is always the user's *conditional* rather than an ending function.

```

23325 \cs_new:Npn \__sort_quick_split:NnNn #1#2#3#4
23326 {
23327   #3 {#2} {#4} \__sort_quick_only_ii:NnnnnNn
23328   \__sort_quick_only_i:NnnnnNn
23329   \__sort_quick_single_end:nnnwnw
23330   { #3 {#4} } { } { } {#2}
23331 }
23332 \cs_new:Npn \__sort_quick_only_i:NnnnnNn #1#2#3#4#5#6#7
23333 {
23334   #6 {#5} {#7} \__sort_quick_split_ii:NnnnnNn
23335   \__sort_quick_only_i:NnnnnNn

```

```

23336     \__sort_quick_only_i_end:nnnwnw
23337     { #6 {#7} } { #3 #2 } { } {#5}
23338 }
23339 \cs_new:Npn \__sort_quick_only_ii:NnnnnNn #1#2#3#4#5#6#7
23340 {
23341     #6 {#5} {#7} \__sort_quick_only_ii:NnnnnNn
23342     \__sort_quick_split_i:NnnnnNn
23343     \__sort_quick_only_i_end:nnnwnw
23344     { #6 {#7} } { } { #4 #2 } {#5}
23345 }
23346 \cs_new:Npn \__sort_quick_split_i:NnnnnNn #1#2#3#4#5#6#7
23347 {
23348     #6 {#5} {#7} \__sort_quick_split_ii:NnnnnNn
23349     \__sort_quick_split_i:NnnnnNn
23350     \__sort_quick_split_end:nnnwnw
23351     { #6 {#7} } { #3 #2 } {#4} {#5}
23352 }
23353 \cs_new:Npn \__sort_quick_split_ii:NnnnnNn #1#2#3#4#5#6#7
23354 {
23355     #6 {#5} {#7} \__sort_quick_split_ii:NnnnnNn
23356     \__sort_quick_split_i:NnnnnNn
23357     \__sort_quick_split_end:nnnwnw
23358     { #6 {#7} } {#3} { #4 #2 } {#5}
23359 }

```

(End definition for __sort_quick_split:NnNn and others.)

```

\__sort_quick_end:nnTFNn
\__sort_quick_single_end:nnnwnw
\__sort_quick_only_i_end:nnnwnw
\__sort_quick_only_ii_end:nnnwnw
\__sort_quick_split_end:nnnwnw

```

The __sort_quick_end:nnTFNn appears instead of the user's conditional, and receives as its arguments the pivot #1, a fake item #2, a `true` and a `false` branches #3 and #4, followed by an ending function #5 (one of the four auxiliaries here) and another copy #6 of the fake item. All those are discarded except the function #5. This function receives lists #1 and #2 of items less than or greater than the pivot #3, then a continuation code #5 just after \s__sort_mark. To avoid a memory problem described earlier, all of the ending functions read #6 until \s__sort_stop and place #6 back into the input stream. When the lists #1 and #2 are empty, the `single` auxiliary simply places the continuation #5 before the pivot {#3}. When #2 is empty, #1 is sorted and placed before the pivot {#3}, taking care to feed the continuation #5 as a continuation for the function sorting #1. When #1 is empty, #2 is sorted, and the continuation argument is used to place the continuation #5 and the pivot {#3} before the sorted result. Finally, when both lists are non-empty, items larger than the pivot are sorted, then items less than the pivot, and the continuations are done in such a way to place the pivot in between.

```

23360 \cs_new:Npn \__sort_quick_end:nnTFNn #1#2#3#4#5#6 {#5}
23361 \cs_new:Npn \__sort_quick_single_end:nnnwnw #1#2#3#4 \s__sort_mark #5#6 \s__sort_stop
23362 { #5 {#3} #6 \s__sort_stop }
23363 \cs_new:Npn \__sort_quick_only_i_end:nnnwnw #1#2#3#4 \s__sort_mark #5#6 \s__sort_stop
23364 {
23365     \__sort_quick_split:NnNn #1
23366     \__sort_quick_end:nnTFNn { } \s__sort_mark {#5}
23367     {#3}
23368     #6 \s__sort_stop
23369 }
23370 \cs_new:Npn \__sort_quick_only_ii_end:nnnwnw #1#2#3#4 \s__sort_mark #5#6 \s__sort_stop
23371 {

```

```

23372     \__sort_quick_split:NnNn #2
23373     \__sort_quick_end:nnTFNn { } \s__sort_mark { #5 {#3} }
23374     #6 \s__sort_stop
23375   }
23376 \cs_new:Npn \__sort_quick_split_end:nnnwnw #1#2#3#4 \s__sort_mark #5#6 \s__sort_stop
23377 {
23378   \__sort_quick_split:NnNn #2 \__sort_quick_end:nnTFNn { } \s__sort_mark
23379   {
23380     \__sort_quick_split:NnNn #1
23381     \__sort_quick_end:nnTFNn { } \s__sort_mark {#5}
23382     {#3}
23383   }
23384   #6 \s__sort_stop
23385 }

```

(End definition for __sort_quick_end:nnTFNn and others.)

39.6 Messages

__sort_error: Bailing out of the sorting code is a bit tricky. It may not be safe to use a delimited argument, so instead we redefine many l3sort commands to be trivial, with __sort_level: jumping to the break point. This error recovery won't work in a group.

```

23386 \cs_new_protected:Npn \__sort_error:
23387 {
23388   \cs_set_eq:NN \__sort_merge_blocks_aux: \prg_do_nothing:
23389   \cs_set_eq:NN \__sort_merge_blocks: \prg_do_nothing:
23390   \cs_set_protected:Npn \__sort_level: { \group_end: \prg_break: }
23391 }

```

(End definition for __sort_error:.)

__sort_disable_toksdef: While sorting, \toksdef is locally disabled to prevent users from using \newtoks or similar commands in their comparison code: the \toks registers that would be assigned are in use by l3sort. In format mode, none of this is needed since there is no \toks allocator.

```

23392 \cs_new_protected:Npn \__sort_disable_toksdef:
23393 { \cs_set_eq:NN \toksdef \__sort_disabled_toksdef:n }
23394 \cs_new_protected:Npn \__sort_disabled_toksdef:n #1
23395 {
23396   \__kernel_msg_error:nnx { kernel } { toksdef }
23397   { \token_to_str:N #1 }
23398   \__sort_error:
23399   \tex_toksdef:D #1
23400 }
23401 \__kernel_msg_new:nnnn { kernel } { toksdef }
23402 { Allocation~of~\iow_char:N\\toks~registers~impossible~while~sorting. }
23403 {
23404   The~comparison~code~used~for~sorting~a~list~has~attempted~to~
23405   define~#1~as~a~new~\iow_char:N\\toks~register~using~
23406   \iow_char:N\\newtoks~
23407   or~a~similar~command.~The~list~will~not~be~sorted.
23408 }

```

(End definition for __sort_disable_toksdef: and __sort_disabled_toksdef:n.)

`__sort_too_long_error:NNw` When there are too many items in a sequence, this is an error, and we clean up properly the mapping over items in the list: break using the type-specific breaking function #1.

```

23409 \cs_new_protected:Npn \__sort_too_long_error:NNw #1#2 \fi:
23410 {
23411   \fi:
23412   \__kernel_msg_error:nnxxx { kernel } { too-large }
23413   { \token_to_str:N #2 }
23414   { \int_eval:n { \l__sort_true_max_int - \l__sort_min_int } }
23415   { \int_eval:n { \l__sort_top_int - \l__sort_min_int } }
23416   #1 \__sort_error:
23417 }
23418 \__kernel_msg_new:nnnn { kernel } { too-large }
23419 { The~list~#1~is~too~long~to~be~sorted~by~TeX. }
23420 {
23421   TeX~has~#2~toks~registers~still~available:~
23422   this~only~allows~to~sort~with~up~to~#3~
23423   items.~The~list~will~not~be~sorted.
23424 }

```

(End definition for `__sort_too_long_error:NNw`.)

```

23425 \__kernel_msg_new:nnnn { kernel } { return-none }
23426 { The~comparison~code~did~not~return. }
23427 {
23428   When~sorting~a~list,~the~code~to~compare~items~#1~and~#2~
23429   did~not~call~
23430   \iow_char:N\sort_return_same: ~nor~
23431   \iow_char:N\sort_return_swapped: .~
23432   Exactly~one~of~these~should~be~called.
23433 }
23434 \__kernel_msg_new:nnnn { kernel } { return-two }
23435 { The~comparison~code~returned~multiple~times. }
23436 {
23437   When~sorting~a~list,~the~code~to~compare~items~#1~and~#2~called~
23438   \iow_char:N\sort_return_same: ~or~
23439   \iow_char:N\sort_return_swapped: ~multiple~times.~
23440   Exactly~one~of~these~should~be~called.
23441 }
23442 </package>

```

40 l3tl-analysis implementation

23443 <@@=tl>

40.1 Internal functions

`\s__tl` The format used to store token lists internally uses the scan mark `\s__tl` as a delimiter.

(End definition for `\s__tl`.)

40.2 Internal format

The task of the l3tl-analysis module is to convert token lists to an internal format which allows us to extract all the relevant information about individual tokens (category code,

character code), as well as reconstruct the token list quickly. This internal format is used in `l3regex` where we need to support arbitrary tokens, and it is used in conversion functions in `l3str-convert`, where we wish to support clusters of characters instead of single tokens.

We thus need a way to encode any $\langle token \rangle$ (even begin-group and end-group character tokens) in a way amenable to manipulating tokens individually. The best we can do is to find $\langle tokens \rangle$ which both *o*-expand and *x*-expand to the given $\langle token \rangle$. Collecting more information about the category code and character code is also useful for regular expressions, since most regexes are catcode-agnostic. The internal format thus takes the form of a succession of items of the form

$\langle tokens \rangle \backslash s_tl \langle catcode \rangle \langle char\ code \rangle \backslash s_tl$

The $\langle tokens \rangle$ *o*- and *x*-expand to the original token in the token list or to the cluster of tokens corresponding to one Unicode character in the given encoding (for `l3str-convert`). The $\langle catcode \rangle$ is given as a single hexadecimal digit, 0 for control sequences. The $\langle char\ code \rangle$ is given as a decimal number, -1 for control sequences.

Using delimited arguments lets us build the $\langle tokens \rangle$ progressively when doing an encoding conversion in `l3str-convert`. On the other hand, the delimiter `\s_tl` may not appear unbraced in $\langle tokens \rangle$. This is not a problem because we are careful to wrap control sequences in braces (as an argument to `\exp_not:n`) when converting from a general token list to the internal format.

The current rule for converting a $\langle token \rangle$ to a balanced set of $\langle tokens \rangle$ which both *o*-expands and *x*-expands to it is the following.

- A control sequence `\cs` becomes `\exp_not:n { \cs } \s_tl 0 -1 \s_tl`.
- A begin-group character `{` becomes `\exp_after:wN { \if_false: } \fi: \s_tl 1 \langle char\ code \rangle \s_tl`.
- An end-group character `}` becomes `\if_false: { \fi: } \s_tl 2 \langle char\ code \rangle \s_tl`.
- A character with any other category code becomes `\exp_not:n { \langle character \rangle } \s_tl \langle hex\ catcode \rangle \langle char\ code \rangle \s_tl`.

23444 `(*package)`

40.3 Variables and helper functions

`\s_tl` The scan mark `\s_tl` is used as a delimiter in the internal format. This is more practical than using a quark, because we would then need to control expansion much more carefully: compare `\int_value:w '#1 \s_tl` with `\int_value:w '#1 \exp_stop_f: \exp_not:N \q_mark` to extract a character code followed by the delimiter in an *x*-expansion.

23445 `\scan_new:N \s_tl`

(End definition for \s_tl.)

`\l_tl_analysis_token` The tokens in the token list are probed with the TeX primitive `\futurelet`. We use `\l_tl_analysis_token` in that construction. In some cases, we convert the following token to a string before probing it: then the token variable used is `\l_tl_analysis_char_token`. When getting tokens from the input stream we may need to look two tokens ahead, for which we use `\l_tl_analysis_next_token`.


```

23446 \cs_new_eq:NN \l__tl_analysis_token ?
23447 \cs_new_eq:NN \l__tl_analysis_char_token ?
23448 \cs_new_eq:NN \l__tl_analysis_next_token ?

```

(End definition for `\l__tl_analysis_token`, `\l__tl_analysis_char_token`, and `\l__tl_analysis_next_token`.)

`\l__tl_peek_code_tl` Holds some code to be run once the next token has been fully analysed in `\peek_analysis_map_inline:n`.

```

23449 \tl_new:N \l__tl_peek_code_tl

```

(End definition for `\l__tl_peek_code_tl`.)

`\c__tl_peek_catcodes_tl` A token list containing the character number 32 (space) with all possible category codes except 1 and 2 (begin-group and end-group). Why 32? Because some LuaTeX versions only allow creation of catcode 10 (space) tokens with this character code, and because even in other engines it is much easier to produce since `\char_generate:nn` refuses to produce spaces.

```

23450 \group_begin:
23451 \char_set_active_eq:NN \ \scan_stop:
23452 \tl_const:Nx \c__tl_peek_catcodes_tl
23453 {
23454   \char_generate:nn { 32 } { 3 } 3
23455   \char_generate:nn { 32 } { 4 } 4
23456   # \char_generate:nn { 32 } { 6 } 6
23457   \char_generate:nn { 32 } { 7 } 7
23458   \char_generate:nn { 32 } { 8 } 8
23459   \c_space_tl \token_to_str:N A
23460   \char_generate:nn { 32 } { 11 } \token_to_str:N B
23461   \char_generate:nn { 32 } { 12 } \token_to_str:N C
23462   \char_generate:nn { 32 } { 13 } \token_to_str:N D
23463 }
23464 \group_end:

```

(End definition for `\c__tl_peek_catcodes_tl`.)

`\l__tl_analysis_normal_int` The number of normal (N-type argument) tokens since the last special token.

```

23465 \int_new:N \l__tl_analysis_normal_int

```

(End definition for `\l__tl_analysis_normal_int`.)

`\l__tl_analysis_index_int` During the first pass, this is the index in the array being built. During the second pass, it is equal to the maximum index in the array from the first pass.

```

23466 \int_new:N \l__tl_analysis_index_int

```

(End definition for `\l__tl_analysis_index_int`.)

`\l__tl_analysis_nesting_int` Nesting depth of explicit begin-group and end-group characters during the first pass. This lets us detect the end of the token list without a reserved end-marker.

```

23467 \int_new:N \l__tl_analysis_nesting_int

```

(End definition for `\l__tl_analysis_nesting_int`.)

`\l__tl_analysis_type_int` When encountering special characters, we record their “type” in this integer.

```

23468 \int_new:N \l__tl_analysis_type_int

```

(End definition for \l__tl_analysis_type_int.)

\g__tl_analysis_result_tl The result of the conversion is stored in this token list, with a succession of items of the form

$\langle tokens \rangle \backslash s_tl \langle catcode \rangle \langle char\ code \rangle \backslash s_tl$

23469 \tl_new:N \g__tl_analysis_result_tl

(End definition for \g__tl_analysis_result_tl.)

_tl_analysis_extract_charcode: Extracting the character code from the meaning of \l__tl_analysis_token. This has no error checking, and should only be assumed to work for begin-group and end-group character tokens. It produces a number in the form ‘ $\langle char \rangle$ ’.

23470 \cs_new:Npn _tl_analysis_extract_charcode:

23471 {

23472 \exp_after:wN _tl_analysis_extract_charcode_aux:w

23473 \token_to_meaning:N \l__tl_analysis_token

23474 }

23475 \cs_new:Npn _tl_analysis_extract_charcode_aux:w #1 ~ #2 ~ { ‘ ’ }

(End definition for _tl_analysis_extract_charcode: and _tl_analysis_extract_charcode_aux:w.)

_tl_analysis_cs_space_count:NN Counts the number of spaces in the string representation of its second argument, as well as the number of characters following the last space in that representation, and feeds the two numbers as semicolon-delimited arguments to the first argument. When this function is used, the escape character is printable and non-space.

23476 \cs_new:Npn _tl_analysis_cs_space_count:NN #1 #2

23477 {

23478 \exp_after:wN #1

23479 \int_value:w \int_eval:w 0

23480 \exp_after:wN _tl_analysis_cs_space_count:w

23481 \token_to_str:N #2

23482 \fi: _tl_analysis_cs_space_count_end:w ; ~ !

23483 }

23484 \cs_new:Npn _tl_analysis_cs_space_count:w #1 ~

23485 {

23486 \if_false: #1 #1 \fi:

23487 + 1

23488 _tl_analysis_cs_space_count:w

23489 }

23490 \cs_new:Npn _tl_analysis_cs_space_count_end:w ; #1 \fi: #2 !

23491 { \exp_after:wN ; \int_value:w \str_count_ignore_spaces:n {#1} ; }

(End definition for _tl_analysis_cs_space_count:NN, _tl_analysis_cs_space_count:w, and _tl_analysis_cs_space_count_end:w.)

40.4 Plan of attack

Our goal is to produce a token list of the form roughly

$\langle token\ 1 \rangle \backslash s@_ \langle catcode\ 1 \rangle \langle char\ code\ 1 \rangle \backslash s@_$
 $\langle token\ 2 \rangle \backslash s_tl \langle catcode\ 2 \rangle \langle char\ code\ 2 \rangle \backslash s_tl$
 $\dots \langle token\ N \rangle \backslash s_tl \langle catcode\ N \rangle \langle char\ code\ N \rangle \backslash s_tl$

Most but not all tokens can be grabbed as an undelimited (N-type) argument by `TEX`. The plan is to have a two pass system. In the first pass, locate special tokens, and store them in various `\toks` registers. In the second pass, which is done within an `x`-expanding assignment, normal tokens are taken in as N-type arguments, and special tokens are retrieved from the `\toks` registers, and removed from the input stream by some means. The whole process takes linear time, because we avoid building the result one item at a time.

We make the escape character printable (backslash, but this later oscillates between slash and backslash): this allows us to distinguish characters from control sequences.

A token has two characteristics: its `\meaning`, and what it looks like for `TEX` when it is in scanning mode (*e.g.*, when capturing parameters for a macro). For our purposes, we distinguish the following meanings:

- begin-group token (category code 1), either space (character code 32), or non-space;
- end-group token (category code 2), either space (character code 32), or non-space;
- space token (category code 10, character code 32);
- anything else (then the token is always an N-type argument).

The token itself can “look like” one of the following

- a non-active character, in which case its meaning is automatically that associated to its character code and category code, we call it “true” character;
- an active character;
- a control sequence.

The only tokens which are not valid N-type arguments are true begin-group characters, true end-group characters, and true spaces. We detect those characters by scanning ahead with `\futurelet`, then distinguishing true characters from control sequences set equal to them using the `\string` representation.

The second pass is a simple exercise in expandable loops.

`__tl_analysis:n` Everything is done within a group, and all definitions are local. We use `\group_align_safe_begin/end:` to avoid problems in case `__tl_analysis:n` is used within an alignment and its argument contains alignment tab tokens.

```

23492 \cs_new_protected:Npn \__tl_analysis:n #1
23493 {
23494   \group_begin:
23495     \group_align_safe_begin:
23496       \__tl_analysis_a:n {#1}
23497       \__tl_analysis_b:n {#1}
23498     \group_align_safe_end:
23499   \group_end:
23500 }
```

(End definition for `__tl_analysis:n`.)

40.5 Disabling active characters

`__tl_analysis_disable:n` Active characters can cause problems later on in the processing, so we provide a way to disable them, by setting them to `undefined`. Since Unicode contains too many characters to loop over all of them, we instead do this whenever we encounter a character. For `pTeX` and `upTeX` we skip characters beyond `[0, 255]` because `\lccode` only allows those values.

```

23501 \group_begin:
23502   \char_set_catcode_active:N \^^@
23503   \cs_new_protected:Npn \__tl_analysis_disable:n #1
23504     {
23505       \tex_lccode:D 0 = #1 \exp_stop_f:
23506       \tex_lowercase:D { \tex_let:D \^^@ } \tex_undefined:D
23507     }
23508   \bool_lazy_or:nnT
23509     { \sys_if_engine_ptex_p: }
23510     { \sys_if_engine_uptex_p: }
23511     {
23512       \cs_gset_protected:Npn \__tl_analysis_disable:n #1
23513         {
23514           \if_int_compare:w 256 > #1 \exp_stop_f:
23515           \tex_lccode:D 0 = #1 \exp_stop_f:
23516           \tex_lowercase:D { \tex_let:D \^^@ } \tex_undefined:D
23517         }
23518     }
23519   }
23520 \group_end:

```

(End definition for `__tl_analysis_disable:n`.)

40.6 First pass

The goal of this pass is to detect special (non-N-type) tokens, and count how many N-type tokens lie between special tokens. Also, we wish to store some representation of each special token in a `\toks` register.

We have 11 types of tokens:

1. a true non-space begin-group character;
2. a true space begin-group character;
3. a true non-space end-group character;
4. a true space end-group character;
5. a true space blank space character;
6. an active character;
7. any other true character;
8. a control sequence equal to a begin-group token (category code 1);
9. a control sequence equal to an end-group token (category code 2);
10. a control sequence equal to a space token (character code 32, category code 10);

11. any other control sequence.

Our first tool is `\futurelet`. This cannot distinguish case 8 from 1 or 2, nor case 9 from 3 or 4, nor case 10 from case 5. Those cases are later distinguished by applying the `\string` primitive to the following token, after possibly changing the escape character to ensure that a control sequence’s string representation cannot be mistaken for the true character.

In cases 6, 7, and 11, the following token is a valid N-type argument, so we grab it and distinguish the case of a character from a control sequence: in the latter case, `\str_tail:n {\token}` is non-empty, because the escape character is printable.

`__tl_analysis_a:n` We read tokens one by one using `\futurelet`. While performing the loop, we keep track of the number of true begin-group characters minus the number of true end-group characters in `\l__tl_analysis_nesting_int`. This reaches -1 when we read the closing brace.

```
23521 \cs_new_protected:Npn \__tl_analysis_a:n #1
23522 {
23523   \__tl_analysis_disable:n { 32 }
23524   \int_set:Nn \tex_escapechar:D { 92 }
23525   \int_zero:N \l__tl_analysis_normal_int
23526   \int_zero:N \l__tl_analysis_index_int
23527   \int_zero:N \l__tl_analysis_nesting_int
23528   \if_false: { \fi: \__tl_analysis_a_loop:w #1 }
23529   \int_decr:N \l__tl_analysis_index_int
23530 }
```

(End definition for `__tl_analysis_a:n`.)

`__tl_analysis_a_loop:w` Read one character and check its type.

```
23531 \cs_new_protected:Npn \__tl_analysis_a_loop:w
23532 { \tex_futurelet:D \l__tl_analysis_token \__tl_analysis_a_type:w }
```

(End definition for `__tl_analysis_a_loop:w`.)

`__tl_analysis_a_type:w` At this point, `\l__tl_analysis_token` holds the meaning of the following token. We store in `\l__tl_analysis_type_int` information about the meaning of the token ahead:

- 0 space token;
- 1 begin-group token;
- -1 end-group token;
- 2 other.

The values 0, 1, -1 correspond to how much a true such character changes the nesting level (2 is used only here, and is irrelevant later). Then call the auxiliary for each case. Note that nesting conditionals here is safe because we only skip over `\l__tl_analysis_token` if it matches with one of the character tokens (hence is not a primitive conditional).

```
23533 \cs_new_protected:Npn \__tl_analysis_a_type:w
23534 {
23535   \l__tl_analysis_type_int =
23536   \if_meaning:w \l__tl_analysis_token \c_space_token
23537     0
```

```

23538     \else:
23539         \if_catcode:w \exp_not:N \l__tl_analysis_token \c_group_begin_token
23540             1
23541     \else:
23542         \if_catcode:w \exp_not:N \l__tl_analysis_token \c_group_end_token
23543             - 1
23544     \else:
23545         2
23546     \fi:
23547 \fi:
23548 \fi:
23549 \exp_stop_f:
23550 \if_case:w \l__tl_analysis_type_int
23551     \exp_after:wN \__tl_analysis_a_space:w
23552 \or: \exp_after:wN \__tl_analysis_a_bgroup:w
23553 \or: \exp_after:wN \__tl_analysis_a_safe:N
23554 \else: \exp_after:wN \__tl_analysis_a_egroup:w
23555 \fi:
23556 }

```

(End definition for __tl_analysis_a_type:w.)

__tl_analysis_a_space:w
 __tl_analysis_a_space_test:w

In this branch, the following token's meaning is a blank space. Apply \string to that token: a true blank space gives a space, a control sequence gives a result starting with the escape character, an active character gives something else than a space since we disabled the space. We grab as \l__tl_analysis_char_token the first character of the string representation then test it in __tl_analysis_a_space_test:w. Also, since __tl_analysis_a_store: expects the special token to be stored in the relevant \toks register, we do that. The extra \exp_not:n is unnecessary of course, but it makes the treatment of all tokens more homogeneous. If we discover that the next token was actually a control sequence or an active character instead of a true space, then we step the counter of normal tokens. We now have in front of us the whole string representation of the control sequence, including potential spaces; those will appear to be true spaces later in this pass. Hence, all other branches of the code in this first pass need to consider the string representation, so that the second pass does not need to test the meaning of tokens, only strings.

```

23557 \cs_new_protected:Npn \__tl_analysis_a_space:w
23558 {
23559     \tex_afterassignment:D \__tl_analysis_a_space_test:w
23560     \exp_after:wN \cs_set_eq:NN
23561     \exp_after:wN \l__tl_analysis_char_token
23562     \token_to_str:N
23563 }
23564 \cs_new_protected:Npn \__tl_analysis_a_space_test:w
23565 {
23566     \if_meaning:w \l__tl_analysis_char_token \c_space_token
23567         \tex_toks:D \l__tl_analysis_index_int { \exp_not:n { ~ } }
23568         \__tl_analysis_a_store:
23569     \else:
23570         \int_incr:N \l__tl_analysis_normal_int
23571     \fi:
23572     \__tl_analysis_a_loop:w
23573 }

```

(End definition for `_tl_analysis_a_space:w` and `_tl_analysis_a_space_test:w`.)

`_tl_analysis_a_bgroup:w` The token is most likely a true character token with catcode 1 or 2, but it might be a control sequence, or an active character. Optimizing for the first case, we store in a toks register some code that expands to that token. Since we will turn what follows into a string, we make sure the escape character is different from the current character code (by switching between solidus and backslash). To detect the special case of an active character let to the catcode 1 or 2 character with the same character code, we disable the active character with that character code and re-test: if the following token has become undefined we can in fact safely grab it. We are finally ready to turn what follows to a string and test it. This is one place where we need `\l_tl_analysis_char_token` to be a separate control sequence from `\l_tl_analysis_token`, to compare them.

```

23574 \group_begin:
23575   \char_set_catcode_group_begin:N \^^@ % {
23576   \cs_new_protected:Npn \_tl\_analysis\_a\_bgroup:w
23577     { \_tl\_analysis\_a\_group:nw { \exp_after:wN \^^@ \if_false: } \fi: } }
23578   \char_set_catcode_group_end:N \^^@
23579   \cs_new_protected:Npn \_tl\_analysis\_a\_egroup:w
23580     { \_tl\_analysis\_a\_group:nw { \if_false: { \fi: \^^@ } } % }
23581 \group_end:
23582 \cs_new_protected:Npn \_tl\_analysis\_a\_group:nw #1
23583 {
23584   \tex_lccode:D 0 = \_tl\_analysis\_extract\_charcode: \scan_stop:
23585   \tex_lowercase:D { \tex_toks:D \l\_tl\_analysis\_index\_int {#1} }
23586   \if_int_compare:w \tex_lccode:D 0 = \tex_escapechar:D
23587     \int_set:Nn \tex_escapechar:D { 139 - \tex_escapechar:D }
23588   \fi:
23589   \_tl\_analysis\_disable:n { \tex_lccode:D 0 }
23590   \tex_futurelet:D \l\_tl\_analysis\_token \_tl\_analysis\_a\_group\_aux:w
23591 }
23592 \cs_new_protected:Npn \_tl\_analysis\_a\_group\_aux:w
23593 {
23594   \if_meaning:w \l\_tl\_analysis\_token \tex_undefined:D
23595     \exp_after:wN \_tl\_analysis\_a\_safe:N
23596   \else:
23597     \exp_after:wN \_tl\_analysis\_a\_group\_auxii:w
23598   \fi:
23599 }
23600 \cs_new_protected:Npn \_tl\_analysis\_a\_group\_auxii:w
23601 {
23602   \tex_afterassignment:D \_tl\_analysis\_a\_group\_test:w
23603   \exp_after:wN \cs_set_eq:NN
23604   \exp_after:wN \l\_tl\_analysis\_char\_token
23605   \token_to_str:N
23606 }
23607 \cs_new_protected:Npn \_tl\_analysis\_a\_group\_test:w
23608 {
23609   \if_charcode:w \l\_tl\_analysis\_token \l\_tl\_analysis\_char\_token
23610     \_tl\_analysis\_a\_store:
23611   \else:
23612     \int_incr:N \l\_tl\_analysis\_normal\_int
23613   \fi:
23614   \_tl\_analysis\_a\_loop:w

```

```
23615 }
```

(End definition for `_tl_analysis_a_bgroup:w` and others.)

`_tl_analysis_a_store:` This function is called each time we meet a special token; at this point, the `\toks` register `\l_tl_analysis_index_int` holds a token list which expands to the given special token. Also, the value of `\l_tl_analysis_type_int` indicates which case we are in:

- -1 end-group character;
- 0 space character;
- 1 begin-group character.

We need to distinguish further the case of a space character (code 32) from other character codes, because those behave differently in the second pass. Namely, after testing the `\lccode` of 0 (which holds the present character code) we change the cases above to

- -2 space end-group character;
- -1 non-space end-group character;
- 0 space blank space character;
- 1 non-space begin-group character;
- 2 space begin-group character.

This has the property that non-space characters correspond to odd values of `\l_tl_analysis_type_int`. The number of normal tokens until here and the type of special token are packed into a `\skip` register. Finally, we check whether we reached the last closing brace, in which case we stop by disabling the looping function (locally).

```
23616 \cs_new_protected:Npn \_tl\_analysis\_a\_store:
23617 {
23618   \tex_advance:D \l\_tl\_analysis\_nesting\_int \l\_tl\_analysis\_type\_int
23619   \if_int_compare:w \tex_lccode:D 0 = '\ \exp_stop_f:
23620     \tex_advance:D \l\_tl\_analysis\_type\_int \l\_tl\_analysis\_type\_int
23621   \fi:
23622   \tex_skip:D \l\_tl\_analysis\_index\_int
23623     = \l\_tl\_analysis\_normal\_int sp
23624     plus \l\_tl\_analysis\_type\_int sp \scan_stop:
23625   \int_incr:N \l\_tl\_analysis\_index\_int
23626   \int_zero:N \l\_tl\_analysis\_normal\_int
23627   \if_int_compare:w \l\_tl\_analysis\_nesting\_int = -1 \exp_stop_f:
23628     \cs_set_eq:NN \_tl\_analysis\_a\_loop:w \scan_stop:
23629   \fi:
23630 }
```

(End definition for `_tl_analysis_a_store:.`)

`_tl_analysis_a_safe:N` This should be the simplest case: since the upcoming token is safe, we can simply grab it in a second pass. If the token is a single character (including space), the `\if_charcode:w` test yields true; we disable a potentially active character (that could otherwise masquerade as the true character in the next pass) and we count one “normal” token. On the other hand, if the token is a control sequence, we should replace it by its string representation for compatibility with other code branches. Instead of slowly looping through

the characters with the main code, we use the knowledge of how the second pass works: if the control sequence name contains no space, count that token as a number of normal tokens equal to its string length. If the control sequence contains spaces, they should be registered as special characters by increasing `_tl_analysis_index_int` (no need to carefully count character between each space), and all characters after the last space should be counted in the following sequence of “normal” tokens.

```

23631 \cs_new_protected:Npn \_tl\_analysis\_a\_safe:N #1
23632 {
23633   \if_charcode:w
23634     \scan_stop:
23635     \exp_after:wN \use_none:n \token_to_str:N #1 \prg_do_nothing:
23636     \scan_stop:
23637     \exp_after:wN \use_i:nn
23638   \else:
23639     \exp_after:wN \use_ii:nn
23640   \fi:
23641   {
23642     \_tl\_analysis\_disable:n { '#1 }
23643     \int_incr:N \_tl\_analysis\_normal\_int
23644   }
23645   { \_tl\_analysis\_cs\_space\_count:NN \_tl\_analysis\_a\_cs:ww #1 }
23646   \_tl\_analysis\_a\_loop:w
23647 }
23648 \cs_new_protected:Npn \_tl\_analysis\_a\_cs:ww #1; #2;
23649 {
23650   \if_int_compare:w #1 > 0 \exp_stop_f:
23651     \tex_skip:D \_tl\_analysis\_index\_int
23652     = \int_eval:n { \_tl\_analysis\_normal\_int + 1 } sp \exp_stop_f:
23653     \tex_advance:D \_tl\_analysis\_index\_int #1 \exp_stop_f:
23654   \else:
23655     \tex_advance:D
23656   \fi:
23657   \_tl\_analysis\_normal\_int #2 \exp_stop_f:
23658 }

```

(End definition for `_tl_analysis_a_safe:N` and `_tl_analysis_a_cs:ww`.)

40.7 Second pass

The second pass is an exercise in expandable loops. All the necessary information is stored in `\skip` and `\toks` registers.

```

\_tl\_analysis\_b:n
\_tl\_analysis\_b\_loop:w

```

Start the loop with the index 0. No need for an end-marker: the loop stops by itself when the last index is read. We repeatedly oscillate between reading long stretches of normal tokens, and reading special tokens.

```

23659 \cs_new_protected:Npn \_tl\_analysis\_b:n #1
23660 {
23661   \_kernel\_tl\_gset:Nx \g\_tl\_analysis\_result\_tl
23662   {
23663     \_tl\_analysis\_b\_loop:w 0; #1
23664     \prg_break_point:
23665   }
23666 }

```

```

23667 \cs_new:Npn \__tl_analysis_b_loop:w #1;
23668 {
23669     \exp_after:wN \__tl_analysis_b_normals:ww
23670     \int_value:w \tex_skip:D #1 ; #1 ;
23671 }

```

(End definition for __tl_analysis_b:n and __tl_analysis_b_loop:w.)

__tl_analysis_b_normals:ww
__tl_analysis_b_normal:wwN

The first argument is the number of normal tokens which remain to be read, and the second argument is the index in the array produced in the first step. A character's string representation is always one character long, while a control sequence is always longer (we have set the escape character to a printable value). In both cases, we leave \exp_not:n {⟨token⟩} \s__tl in the input stream (after x-expansion). Here, \exp_not:n is used rather than \exp_not:N because #3 could be a macro parameter character or could be \s__tl (which must be hidden behind braces in the result).

```

23672 \cs_new:Npn \__tl_analysis_b_normals:ww #1;
23673 {
23674     \if_int_compare:w #1 = 0 \exp_stop_f:
23675     \__tl_analysis_b_special:w
23676     \fi:
23677     \__tl_analysis_b_normal:wwN #1;
23678 }
23679 \cs_new:Npn \__tl_analysis_b_normal:wwN #1; #2; #3
23680 {
23681     \exp_not:n { \exp_not:n { #3 } } \s__tl
23682     \if_charcode:w
23683         \scan_stop:
23684         \exp_after:wN \use_none:n \token_to_str:N #3 \prg_do_nothing:
23685         \scan_stop:
23686         \exp_after:wN \__tl_analysis_b_char:Nww
23687     \else:
23688         \exp_after:wN \__tl_analysis_b_cs:Nww
23689     \fi:
23690     #3 #1; #2;
23691 }

```

(End definition for __tl_analysis_b_normals:ww and __tl_analysis_b_normal:wwN.)

__tl_analysis_b_char:Nww

If the normal token we grab is a character, leave ⟨catcode⟩ ⟨charcode⟩ followed by \s__tl in the input stream, and call __tl_analysis_b_normals:ww with its first argument decremented.

```

23692 \cs_new:Npx \__tl_analysis_b_char:Nww #1
23693 {
23694     \exp_not:N \if_meaning:w #1 \exp_not:N \tex_undefined:D
23695     \token_to_str:N D \exp_not:N \else:
23696     \exp_not:N \if_catcode:w #1 \c_catcode_other_token
23697     \token_to_str:N C \exp_not:N \else:
23698     \exp_not:N \if_catcode:w #1 \c_catcode_letter_token
23699     \token_to_str:N B \exp_not:N \else:
23700     \exp_not:N \if_catcode:w #1 \c_math_toggle_token      3
23701     \exp_not:N \else:
23702     \exp_not:N \if_catcode:w #1 \c_alignment_token      4
23703     \exp_not:N \else:
23704     \exp_not:N \if_catcode:w #1 \c_math_superscript_token 7

```

```

23705     \exp_not:N \else:
23706 \exp_not:N \if_catcode:w #1 \c_math_subscript_token 8
23707     \exp_not:N \else:
23708 \exp_not:N \if_catcode:w #1 \c_space_token
23709     \token_to_str:N A \exp_not:N \else:
23710     6
23711 \exp_not:n { \fi: \fi: \fi: \fi: \fi: \fi: \fi: \fi: }
23712 \exp_not:N \int_value:w '#1 \s__tl
23713 \exp_not:N \exp_after:wN \exp_not:N \__tl_analysis_b_normals:ww
23714 \exp_not:N \int_value:w \exp_not:N \int_eval:w - 1 +
23715 }

```

(End definition for __tl_analysis_b_char:Nww.)

__tl_analysis_b_cs:Nww If the token we grab is a control sequence, leave 0 -1 (as category code and character code) in the input stream, followed by \s__tl, and call __tl_analysis_b_normals:ww with updated arguments.

```

23716 \cs_new:Npn \__tl_analysis_b_cs:Nww #1
23717 {
23718     0 -1 \s__tl
23719     \__tl_analysis_cs_space_count:NN \__tl_analysis_b_cs_test:ww #1
23720 }
23721 \cs_new:Npn \__tl_analysis_b_cs_test:ww #1 ; #2 ; #3 ; #4 ;
23722 {
23723     \exp_after:wN \__tl_analysis_b_normals:ww
23724     \int_value:w \int_eval:w
23725     \if_int_compare:w #1 = 0 \exp_stop_f:
23726     #3
23727     \else:
23728     \tex_skip:D \int_eval:n { #4 + #1 } \exp_stop_f:
23729     \fi:
23730     - #2
23731     \exp_after:wN ;
23732     \int_value:w \int_eval:n { #4 + #1 } ;
23733 }

```

(End definition for __tl_analysis_b_cs:Nww and __tl_analysis_b_cs_test:ww.)

__tl_analysis_b_special:w Here, #1 is the current index in the array built in the first pass. Check now whether we reached the end (we shouldn't keep the trailing end-group character that marked the end of the token list in the first pass). Unpack the \toks register: when x-expanding again, we will get the special token. Then leave the category code in the input stream, followed by the character code, and call __tl_analysis_b_loop:w with the next index.

```

23734 \group_begin:
23735     \char_set_catcode_other:N A
23736     \cs_new:Npn \__tl_analysis_b_special:w
23737     \fi: \__tl_analysis_b_normal:wwN 0 ; #1 ;
23738     {
23739     \fi:
23740     \if_int_compare:w #1 = \l__tl_analysis_index_int
23741     \exp_after:wN \prg_break:
23742     \fi:
23743     \tex_the:D \tex_toks:D #1 \s__tl
23744     \if_case:w \tex_gluestretch:D \tex_skip:D #1 \exp_stop_f:

```

```

23745         \token_to_str:N A
23746     \or: 1
23747     \or: 1
23748     \else: 2
23749     \fi:
23750     \if_int_odd:w \tex_gluestretch:D \tex_skip:D #1 \exp_stop_f:
23751         \exp_after:wN \_tl_analysis_b_special_char:wN \int_value:w
23752     \else:
23753         \exp_after:wN \_tl_analysis_b_special_space:w \int_value:w
23754     \fi:
23755     \int_eval:n { 1 + #1 } \exp_after:wN ;
23756     \token_to_str:N
23757 }
23758 \group_end:
23759 \cs_new:Npn \_tl_analysis_b_special_char:wN #1 ; #2
23760 {
23761     \int_value:w '#2 \s_tl
23762     \_tl_analysis_b_loop:w #1 ;
23763 }
23764 \cs_new:Npn \_tl_analysis_b_special_space:w #1 ; ~
23765 {
23766     32 \s_tl
23767     \_tl_analysis_b_loop:w #1 ;
23768 }

```

(End definition for `_tl_analysis_b_special:w`, `_tl_analysis_b_special_char:wN`, and `_tl_analysis_b_special_space:w`.)

40.8 Mapping through the analysis

```

\tl_analysis_map_inline:nn
\tl_analysis_map_inline:Nn
  \_tl_analysis_map_inline_aux:Nn
  \_tl_analysis_map_inline_aux:nnn

```

First obtain the analysis of the token list into `\g__tl_analysis_result_tl`. To allow nested mappings, increase the nesting depth `\g__kernel_prg_map_int` (shared between all modules), then define the looping macro, which has a name specific to that nesting depth. That looping grabs the `<tokens>`, `<catcode>` and `<char code>`; it checks for the end of the loop with `\use_none:n ##2`, normally empty, but which becomes `\tl_map_break:` at the end; it then performs the user's code `#2`, and loops by calling itself. When the loop ends, remember to decrease the nesting depth.

```

23769 \cs_new_protected:Npn \tl_analysis_map_inline:nn #1
23770 {
23771     \_tl_analysis:n {#1}
23772     \int_gincr:N \g__kernel_prg_map_int
23773     \exp_args:Nc \_tl_analysis_map_inline_aux:Nn
23774         { \_tl_analysis_map_inline_ \int_use:N \g__kernel_prg_map_int :wNw }
23775 }
23776 \cs_new_protected:Npn \tl_analysis_map_inline:Nn #1
23777 { \exp_args:No \tl_analysis_map_inline:nn #1 }
23778 \cs_new_protected:Npn \_tl_analysis_map_inline_aux:Nn #1#2
23779 {
23780     \cs_gset_protected:Npn #1 ##1 \s_tl ##2 ##3 \s_tl
23781     {
23782         \use_none:n ##2
23783         \_tl_analysis_map_inline_aux:nnn {##1} {##3} {##2}
23784     }
23785     \cs_gset_protected:Npn \_tl_analysis_map_inline_aux:nnn ##1##2##3

```

```

23786     {
23787         #2
23788         #1
23789     }
23790     \exp_after:wN #1
23791     \g__tl_analysis_result_tl
23792     \s__tl { ? \tl_map_break: } \s__tl
23793     \prg_break_point:Nn \tl_map_break:
23794     { \int_gdecr:N \g__kernel_prg_map_int }
23795 }

```

(End definition for `\tl_analysis_map_inline:nn` and others. These functions are documented on page 228.)

40.9 Showing the results

`\tl_analysis_show:N` Add to `__tl_analysis:n` a third pass to display tokens to the terminal. If the token list variable is not defined, throw the same error as `\tl_show:N` by simply calling that function.

```

23796 \cs_new_protected:Npn \tl_analysis_show:N #1
23797 {
23798     \tl_if_exist:NTF #1
23799     {
23800         \exp_args:No \__tl_analysis:n {#1}
23801         \msg_show:nnxxxx { LaTeX / kernel } { show-tl-analysis }
23802         { \token_to_str:N #1 } { \__tl_analysis_show: } { } { }
23803     }
23804     { \tl_show:N #1 }
23805 }
23806 \cs_new_protected:Npn \tl_analysis_show:n #1
23807 {
23808     \__tl_analysis:n {#1}
23809     \msg_show:nnxxxx { LaTeX / kernel } { show-tl-analysis }
23810     { } { \__tl_analysis_show: } { } { }
23811 }

```

(End definition for `\tl_analysis_show:N` and `\tl_analysis_show:n`. These functions are documented on page 228.)

`__tl_analysis_show:` Here, `#1` o- and x-expands to the token; `#2` is the category code (one uppercase hexadecimal digit), 0 for control sequences; `#3` is the character code, which we ignore. In the cases of control sequences and active characters, the meaning may overflow one line, and we want to truncate it. Those cases are thus separated out.

```

23812 \cs_new:Npn \__tl_analysis_show:
23813 {
23814     \exp_after:wN \__tl_analysis_show_loop:wNw \g__tl_analysis_result_tl
23815     \s__tl { ? \prg_break: } \s__tl
23816     \prg_break_point:
23817 }
23818 \cs_new:Npn \__tl_analysis_show_loop:wNw #1 \s__tl #2 #3 \s__tl
23819 {
23820     \use_none:n #2
23821     \iow_newline: > \use:nn { ~ } { ~ }
23822     \if_int_compare:w "#2 = 0 \exp_stop_f:

```

```

23823     \exp_after:wN \__tl_analysis_show_cs:n
23824 \else:
23825     \if_int_compare:w "#2 = 13 \exp_stop_f:
23826         \exp_after:wN \exp_after:wN
23827         \exp_after:wN \__tl_analysis_show_active:n
23828     \else:
23829         \exp_after:wN \exp_after:wN
23830         \exp_after:wN \__tl_analysis_show_normal:n
23831     \fi:
23832 \fi:
23833 {#1}
23834 \__tl_analysis_show_loop:wNw
23835 }

```

(End definition for __tl_analysis_show: and __tl_analysis_show_loop:wNw.)

__tl_analysis_show_normal:n Non-active characters are a simple matter of printing the character, and its meaning. Our test suite checks that begin-group and end-group characters do not mess up T_EX's alignment status.

```

23836 \cs_new:Npn \__tl_analysis_show_normal:n #1
23837 {
23838     \exp_after:wN \token_to_str:N #1 ~
23839     ( \exp_after:wN \token_to_meaning:N #1 )
23840 }

```

(End definition for __tl_analysis_show_normal:n.)

__tl_analysis_show_value:N This expands to the value of #1 if it has any.

```

23841 \cs_new:Npn \__tl_analysis_show_value:N #1
23842 {
23843     \token_if_expandable:NF #1
23844     {
23845         \token_if_chardef:NTF #1 \prg_break: { }
23846         \token_if_mathchardef:NTF #1 \prg_break: { }
23847         \token_if_dim_register:NTF #1 \prg_break: { }
23848         \token_if_int_register:NTF #1 \prg_break: { }
23849         \token_if_skip_register:NTF #1 \prg_break: { }
23850         \token_if_toks_register:NTF #1 \prg_break: { }
23851         \use_none:nnn
23852         \prg_break_point:
23853         \use:n { \exp_after:wN = \tex_the:D #1 }
23854     }
23855 }

```

(End definition for __tl_analysis_show_value:N.)

__tl_analysis_show_cs:n Control sequences and active characters are printed in the same way, making sure not to go beyond the \l_iow_line_count_int. In case of an overflow, we replace the last characters by \c__tl_analysis_show_etc_str.

```

\__tl_analysis_show_active:n
\__tl_analysis_show_long:nn
\__tl_analysis_show_long_aux:nnnn
23856 \cs_new:Npn \__tl_analysis_show_cs:n #1
23857 { \exp_args:No \__tl_analysis_show_long:nn {#1} { control~sequence= } }
23858 \cs_new:Npn \__tl_analysis_show_active:n #1
23859 { \exp_args:No \__tl_analysis_show_long:nn {#1} { active~character= } }
23860 \cs_new:Npn \__tl_analysis_show_long:nn #1

```

```

23861 {
23862   \_tl_analysis_show_long_aux:oofn
23863   { \token_to_str:N #1 }
23864   { \token_to_meaning:N #1 }
23865   { \_tl_analysis_show_value:N #1 }
23866 }
23867 \cs_new:Npn \_tl_analysis_show_long_aux:nnnn #1#2#3#4
23868 {
23869   \int_compare:nNnTF
23870     { \str_count:n { #1 ~ ( #4 #2 #3 ) } }
23871     > { \l_iow_line_count_int - 3 }
23872     {
23873       \str_range:nnn { #1 ~ ( #4 #2 #3 ) } { 1 }
23874       {
23875         \l_iow_line_count_int - 3
23876         - \str_count:N \c__tl_analysis_show_etc_str
23877       }
23878       \c__tl_analysis_show_etc_str
23879     }
23880     { #1 ~ ( #4 #2 #3 ) }
23881   }
23882 \cs_generate_variant:Nn \_tl_analysis_show_long_aux:nnnn { oof }

```

(End definition for `_tl_analysis_show_cs:n` and others.)

40.10 Peeking ahead

`\peek_analysis_map_break:` The break statements use the general `\prg_map_break:Nn`.
`\peek_analysis_map_break:n`

```

23883 \cs_new:Npn \peek_analysis_map_break:
23884   { \prg_map_break:Nn \peek_analysis_map_break: { } }
23885 \cs_new:Npn \peek_analysis_map_break:n
23886   { \prg_map_break:Nn \peek_analysis_map_break: }

```

(End definition for `\peek_analysis_map_break:` and `\peek_analysis_map_break:n`. These functions are documented on page [142](#).)

`\l__tl_peek_charcode_int`

```

23887 \int_new:N \l__tl_peek_charcode_int

```

(End definition for `\l__tl_peek_charcode_int`.)

`_tl_analysis_char_arg:Nw`
`_tl_analysis_char_arg_aux:Nw`

After a call to `\futurelet \l__tl_analysis_token` followed by a stringified character token (either explicit space or catcode other character), grab the argument and pass it to #1. We only need to do anything in the case of a space.

```

23888 \cs_new:Npn \_tl_analysis_char_arg:Nw
23889 {
23890   \if_meaning:w \l__tl_analysis_token \c_space_token
23891     \exp_after:wN \_tl_analysis_char_arg_aux:Nw
23892   \fi:
23893 }
23894 \cs_new:Npn \_tl_analysis_char_arg_aux:Nw #1 ~ { #1 { ~ } }

```

(End definition for `_tl_analysis_char_arg:Nw` and `_tl_analysis_char_arg_aux:Nw`.)

```

\peek_analysis_map_inline:n
\__tl_peek_analysis_loop:NNn
\__tl_peek_analysis_test:
\__tl_peek_analysis_normal:N
\__tl_peek_analysis_cs:
\__tl_peek_analysis_char:N
\__tl_peek_analysis_char:nN
\__tl_peek_analysis_special:
\__tl_peek_analysis_retest:
\__tl_peek_analysis_next:
\__tl_peek_analysis_str:
\__tl_peek_analysis_str:w
\__tl_peek_analysis_str:n
\__tl_peek_analysis_active_str:n
\__tl_peek_analysis_explicit:n
\__tl_peek_analysis_escape:
\__tl_peek_analysis_collect:w
\__tl_peek_analysis_collect:n
\__tl_peek_analysis_collect_loop:
\__tl_peek_analysis_collect_test:
\__tl_peek_analysis_collect_end:NNN

```

Save the user's code in a control sequence that is suitable for nested maps. We may wish to pass to this function an `\outer` control sequence or active character; for this we will undefine potentially-`\outer` tokens within a group, closed after the function receives its arguments. This user's code function also calls the loop auxiliary, and includes the trailing `\prg_break_point:Nn` for when the user wants to stop the loop. The loop auxiliary must remove that break point because it must look at the input stream.

```

23895 \cs_new_protected:Npn \peek_analysis_map_inline:n #1
23896 {
23897   \int_gincr:N \g__kernel_prg_map_int
23898   \cs_set_protected:cpn
23899   { __tl_analysis_map_ \int_use:N \g__kernel_prg_map_int :nnN }
23900   ##1##2##3
23901   {
23902     \group_end:
23903     #1
23904     \__tl_peek_analysis_loop:NNn
23905     \prg_break_point:Nn \peek_analysis_map_break: { }
23906   }
23907   \__tl_peek_analysis_loop:NNn ? ? ?
23908 }

```

The loop starts a group (closed by the user-code function defined above) with a normalized escape character, and checks if the next token is special or N-type.

```

23909 \cs_new_protected:Npn \__tl_peek_analysis_loop:NNn #1#2#3
23910 {
23911   \group_begin:
23912   \tl_set:Nx \l__tl_peek_code_tl
23913   {
23914     \exp_not:c
23915     { __tl_analysis_map_ \int_use:N \g__kernel_prg_map_int :nnN }
23916   }
23917   \int_set:Nn \tex_escapechar:D { '\ }
23918   \peek_after:Nw \__tl_peek_analysis_test:
23919 }
23920 \cs_new_protected:Npn \__tl_peek_analysis_test:
23921 {
23922   \if_int_odd:w
23923   \if_catcode:w \exp_not:N \l_peek_token { 0 \exp_stop_f: \fi:
23924   \if_catcode:w \exp_not:N \l_peek_token } 0 \exp_stop_f: \fi:
23925   \if_meaning:w \l_peek_token \c_space_token 0 \exp_stop_f: \fi:
23926   1 \exp_stop_f:
23927   \exp_after:wN \exp_after:wN
23928   \exp_after:wN \__tl_peek_analysis_normal:N
23929   \exp_after:wN \exp_not:N
23930   \else:
23931   \exp_after:wN \__tl_peek_analysis_special:
23932   \fi:
23933 }

```

Normal tokens are not too hard, but can be `\outer`, hence the `\exp_not:N` in the code above. If the token is expandable then it might be an `\outer` or a TeX conditional, so to be safe we set it to `\scan_stop:` (the assignment is local and stopped by the `\group_end:` upon calling the user's code). Then distinguish characters (including active ones and macro parameter characters) from control sequences (whose string representation

is more than one character because the escape character is printable). For a control sequence call the user code with suitable arguments.

```

23934 \cs_new_protected:Npn \__tl_peek_analysis_normal:N #1
23935 {
23936   \exp_after:wN \reverse_if:N \exp_after:wN \if_meaning:w
23937   \exp_not:N #1 #1
23938   \tex_let:D #1 \scan_stop:
23939   \tl_put_right:Nn \l__tl_peek_code_tl { { \exp_not:N #1 } }
23940   \else:
23941     \tl_put_right:Nn \l__tl_peek_code_tl { { \exp_not:n {#1} } }
23942   \fi:
23943   \if_charcode:w
23944     \scan_stop:
23945     \exp_after:wN \use_none:n \token_to_str:N #1 \prg_do_nothing:
23946     \scan_stop:
23947     \exp_after:wN \__tl_peek_analysis_char:N
23948     \exp_after:wN #1
23949   \else:
23950     \exp_after:wN \__tl_peek_analysis_cs:
23951   \fi:
23952 }
23953 \cs_new_protected:Npn \__tl_peek_analysis_cs:
23954 { \l__tl_peek_code_tl { -1 } 0 }
23955 \cs_new_protected:Npn \__tl_peek_analysis_char:N #1
23956 {
23957   \char_set_lccode:nn { '#1 } { 32 }
23958   \tex_lowercase:D { \__tl_peek_analysis_char:nN {#1} } #1
23959 }
23960 \cs_new_protected:Npn \__tl_peek_analysis_char:nN #1#2
23961 {
23962   \cs_set_protected:Npn \__tl_tmp:w ##1 #1 ##2 ##3 \scan_stop:
23963   { \exp_args:No \l__tl_peek_code_tl { \int_value:w '#2 } ##2 }
23964   \exp_after:wN \__tl_tmp:w \c__tl_peek_catcodes_tl \scan_stop:
23965 }

```

For special characters the idea is to eventually act with `\token_to_str:N`, then pick up one by one the characters of this string representation until hitting the token that follows. First determine the character code of (the meaning of) the *<token>* (which we know is a special token), make sure the escape character is different from it, normalize the meanings of two active characters and the empty control sequence, and filter out these cases in `__tl_peek_analysis_retest:`.

```

23966 \cs_new_protected:Npn \__tl_peek_analysis_special:
23967 {
23968   \tex_let:D \l__tl_analysis_token = ~ \l_peek_token
23969   \int_set:Nn \l__tl_peek_charcode_int
23970   { \__tl_analysis_extract_charcode: }
23971   \if_int_compare:w \l__tl_peek_charcode_int = \tex_escapechar:D
23972   \int_set:Nn \tex_escapechar:D { '\ / }
23973   \fi:
23974   \char_set_active_eq:nN { \l__tl_peek_charcode_int } \scan_stop:
23975   \char_set_active_eq:nN { \tex_escapechar:D } \scan_stop:
23976   \cs_set_eq:cN { } \scan_stop:
23977   \tex_futurelet:D \l__tl_analysis_token
23978   \__tl_peek_analysis_retest:

```

```

23979 }
23980 \cs_new_protected:Npn \__tl_peek_analysis_retest:
23981 {
23982   \if_meaning:w \l__tl_analysis_token \scan_stop:
23983     \exp_after:wN \__tl_peek_analysis_normal:N
23984   \else:
23985     \exp_after:wN \__tl_peek_analysis_next:
23986   \fi:
23987 }

```

At this point we know the meaning of the $\langle token \rangle$ in the input stream is $\backslash l_peek_token$, either a space (32, 10) or a begin-group or end-group token (catcode 1 or 2), and we excluded a few cases that would be difficult later (empty control sequence, active character with the same character code as its meaning or as the escape character). Now look at the $\langle next token \rangle$ following it using a combination of $\backslash afterassignment$ and $\backslash futurelet$. The syntax of this primitive is $\backslash futurelet \langle peek token \rangle \langle first token \rangle \langle next token \rangle$, and it sets $\langle peek token \rangle$ equal to $\langle next token \rangle$. Traditionally, one takes $\langle first token \rangle$ to be some macro that regains control of the code and, e.g., analyses $\langle peek token \rangle$. Here, both $\langle first token \rangle$ and $\langle next token \rangle$ are mostly unknown tokens in the input stream (but we know the $\langle first token \rangle$ has catcode 1, 2 or 10), where $\langle first token \rangle$ was already stored as $\backslash l_peek_token$, and we regain control using $\backslash afterassignment$, which inserts its argument after the assignment, hence after $\langle peek token \rangle$ but before $\langle first token \rangle$.

```

23988 \cs_new_protected:Npn \__tl_peek_analysis_next:
23989 {
23990   \tl_if_empty:oT { \tex_the:D \tex_everyeof:D }
23991   { \tex_everyeof:D { \scan_stop: } }
23992   \tex_afterassignment:D \__tl_peek_analysis_str:
23993   \tex_futurelet:D \l__tl_analysis_next_token
23994 }

```

We then hit the $\langle first token \rangle$ with $\backslash token_to_str:N$ and grab characters until finding $\backslash l_tl_analysis_next_token$. More precisely, by looking at the first character in the string representation of the $\langle first token \rangle$ we distinguish three cases: a stringified control sequence starts with the escape character; for an explicit character we find that same character; for an explicit character we find anything else (we made sure to exclude the case of an active character whose string representation coincides with the other two cases).

```

23995 \cs_new_protected:Npn \__tl_peek_analysis_str:
23996 {
23997   \exp_after:wN \tex_futurelet:D
23998   \exp_after:wN \l__tl_analysis_token
23999   \exp_after:wN \__tl_peek_analysis_str:w
24000   \token_to_str:N
24001 }
24002 \cs_new_protected:Npn \__tl_peek_analysis_str:w
24003 { \__tl_analysis_char_arg:Nw \__tl_peek_analysis_str:n }
24004 \cs_new_protected:Npn \__tl_peek_analysis_str:n #1
24005 {
24006   \int_case:nnF { '#1 }
24007   {
24008     { \l__tl_peek_charcode_int }
24009     { \__tl_peek_analysis_explicit:n {#1} }
24010     { \tex_escapechar:D } { \__tl_peek_analysis_escape: }
24011   }

```

```

24012     { \_tl_peek_analysis_active_str:n {#1} }
24013   }

```

When #1 is a stringified active character we pass appropriate arguments to the user's code; thankfully `\char_generate:nn` can make active characters.

```

24014 \cs_new_protected:Npn \_tl_peek_analysis_active_str:n #1
24015 {
24016   \tl_put_right:Nx \l__tl_peek_code_tl
24017   {
24018     { \char_generate:nn { '#1 } { 13 } }
24019     { \int_value:w '#1 }
24020     \token_to_str:N D
24021   }
24022   \l__tl_peek_code_tl
24023 }

```

When #1 matches the character we had extracted from the meaning of `\l_peek_token`, the token was an explicit character, which can be a standard space, or a begin-group or end-group character with some character code. In the latter two cases we call `\char_generate:nn` with suitable arguments and put suitable `\if_false: \fi:` constructions to make the result balanced and such that o-expanding or x-expanding gives back a single (unbalanced) begin-group or end-group character.

```

24024 \cs_new_protected:Npn \_tl_peek_analysis_explicit:n #1
24025 {
24026   \tl_put_right:Nx \l__tl_peek_code_tl
24027   {
24028     \if_meaning:w \l_peek_token \c_space_token
24029     { ~ } { 32 } \token_to_str:N A
24030     \else:
24031       \if_catcode:w \l_peek_token \c_group_begin_token
24032       {
24033         \exp_not:N \exp_after:wN
24034         \char_generate:nn { '#1 } { 1 }
24035         \exp_not:N \if_false:
24036         \if_false: { \fi: }
24037         \exp_not:N \fi:
24038       }
24039       { \int_value:w '#1 }
24040       1
24041     \else:
24042     {
24043       \exp_not:N \if_false:
24044       { \if_false: } \fi:
24045       \exp_not:N \fi:
24046       \char_generate:nn { '#1 } { 2 }
24047     }
24048     { \int_value:w '#1 }
24049     2
24050     \fi:
24051   }
24052   \l__tl_peek_code_tl
24053 }
24054

```

Finally there is the case of a special token whose string representation starts with an escape character, namely the token was a control sequence. In that case we could have grabbed the token directly as an N-type argument, but of course we couldn't know that until we had run all the various tests including stringifying the token. We are thus left with the hard work of picking up one by one the characters in the csname (being careful about spaces), until finding a token that matches the *<next token>* picked up earlier (which was not stringified), such that the control sequence that we found so far indeed has the expected meaning `\l_peek_token`. This comparison with `\l_peek_token` catches a reasonably common case like `\c_group_begin_token _` in which the trailing `_` has category code other: without comparison of the constructed csname with `\l_peek_token` collection would stop at `\c`, which is wrong.

```

24055 \cs_new_protected:Npn \__tl_peek_analysis_escape:
24056 {
24057   \tl_clear:N \l__tl_internal_a_tl
24058   \tex_futurelet:D \l__tl_analysis_token
24059   \__tl_peek_analysis_collect:w
24060 }
24061 \cs_new_protected:Npn \__tl_peek_analysis_collect:w
24062 { \__tl_analysis_char_arg:Nw \__tl_peek_analysis_collect:n }
24063 \cs_new_protected:Npn \__tl_peek_analysis_collect:n #1
24064 {
24065   \tl_put_right:Nn \l__tl_internal_a_tl {#1}
24066   \__tl_peek_analysis_collect_loop:
24067 }
24068 \cs_new_protected:Npn \__tl_peek_analysis_collect_loop:
24069 {
24070   \tex_futurelet:D \l__tl_analysis_token
24071   \__tl_peek_analysis_collect_test:
24072 }
24073 \cs_new_protected:Npn \__tl_peek_analysis_collect_test:
24074 {
24075   \if_meaning:w \l__tl_analysis_token \l__tl_analysis_next_token
24076   \exp_after:wN \if_meaning:w \cs:w \l__tl_internal_a_tl \cs_end: \l_peek_token
24077   \__tl_peek_analysis_collect_end:NNN
24078   \fi:
24079   \fi:
24080   \__tl_peek_analysis_collect:w
24081 }

```

End by calling the user code with suitable arguments (here #1, #2 are `\fi:`), which closes the group begun early on.

```

24082 \cs_new_protected:Npn \__tl_peek_analysis_collect_end:NNN #1#2#3
24083 {
24084   #1 #2
24085   \tl_put_right:Nx \l__tl_peek_code_tl
24086   {
24087     { \exp_not:N \exp_not:n { \exp_not:c { \l__tl_internal_a_tl } } }
24088     { -1 }
24089     0
24090   }
24091   \l__tl_peek_code_tl
24092 }

```

(End definition for `\peek_analysis_map_inline:n` and others. This function is documented on page 142.)

40.11 Messages

`\c__tl_analysis_show_etc_str` When a control sequence (or active character) and its meaning are too long to fit in one line of the terminal, the end is replaced by this token list.

```
24093 \tl_const:Nx \c__tl_analysis_show_etc_str % (
24094   { \token_to_str:N \ETC.) }
```

(End definition for `\c__tl_analysis_show_etc_str`.)

```
24095 \__kernel_msg_new:nnn { kernel } { show-tl-analysis }
24096   {
24097     The-token-list~ \tl_if_empty:nF {#1} { #1 ~ }
24098     \tl_if_empty:nTF {#2}
24099       { is-empty }
24100       { contains-the-tokens: #2 }
24101   }
24102 </package>
```

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```
24103 <*package>
24104 <@@=regex>
```

41.1 Plan of attack

Most regex engines use backtracking. This allows to provide very powerful features (back-references come to mind first), but it is costly, and raises the problem of catastrophic backtracking. Since \TeX is not first and foremost a programming language, complicated code tends to run slowly, and we must use faster, albeit slightly more restrictive, techniques, coming from automata theory.

Given a regular expression of n characters, we do the following:

- (Compiling.) Analyse the regex, finding invalid input, and convert it to an internal representation.
- (Building.) Convert the compiled regex to a non-deterministic finite automaton (NFA) with $O(n)$ states which accepts precisely token lists matching that regex.
- (Matching.) Loop through the query token list one token (one “position”) at a time, exploring in parallel every possible path (“active thread”) through the NFA, considering active threads in an order determined by the quantifiers’ greediness.

We use the following vocabulary in the code comments (and in variable names).

- *Group*: index of the capturing group, -1 for non-capturing groups.
- *Position*: each token in the query is labelled by an integer $\langle position \rangle$, with $\text{min_pos} - 1 \leq \langle position \rangle \leq \text{max_pos}$. The lowest and highest positions $\text{min_pos} - 1$ and max_pos correspond to imaginary begin and end markers (with non-existent category code and character code). max_pos is only set quite late in the processing.

- *Query*: the token list to which we apply the regular expression.
- *State*: each state of the NFA is labelled by an integer $\langle state \rangle$ with $\text{min_state} \leq \langle state \rangle < \text{max_state}$.
- *Active thread*: state of the NFA that is reached when reading the query token list for the matching. Those threads are ordered according to the greediness of quantifiers.
- *Step*: used when matching, starts at 0, incremented every time a character is read, and is not reset when searching for repeated matches. The integer `\l__regex_step_int` is a unique id for all the steps of the matching algorithm.

We use `l3intarray` to manipulate arrays of integers. We also abuse \TeX 's `\toks` registers, by accessing them directly by number rather than tying them to control sequence using the `\newtoks` allocation functions. Specifically, these arrays and `\toks` are used as follows. When building, `\toks $\langle state \rangle$` holds the tests and actions to perform in the $\langle state \rangle$ of the NFA. When matching,

- `\g__regex_state_active_intarray` holds the last $\langle step \rangle$ in which each $\langle state \rangle$ was active.
- `\g__regex_thread_info_intarray` consists of blocks for each $\langle thread \rangle$ (with $\text{min_thread} \leq \langle thread \rangle < \text{max_thread}$). Each block has $1+2\text{\l__regex_capturing_group_int}$ entries: the $\langle state \rangle$ in which the $\langle thread \rangle$ currently is, followed by the beginnings of all submatches, and then the ends of all submatches. The $\langle threads \rangle$ are ordered starting from the best to the least preferred.
- `\g__regex_submatch_prev_intarray`, `\g__regex_submatch_begin_intarray` and `\g__regex_submatch_end_intarray` hold, for each submatch (as would be extracted by `\regex_extract_all:nnN`), the place where the submatch started to be looked for and its two end-points. For historical reasons, the minimum index is twice max_state , and the used registers go up to `\l__regex_submatch_int`. They are organized in blocks of `\l__regex_capturing_group_int` entries, each block corresponding to one match with all its submatches stored in consecutive entries.

When actually building the result,

- `\toks $\langle position \rangle$` holds $\langle tokens \rangle$ which o- and x-expand to the $\langle position \rangle$ -th token in the query.
- `\g__regex_balance_intarray` holds the balance of begin-group and end-group character tokens which appear before that point in the token list.

The code is structured as follows. Variables are introduced in the relevant section. First we present some generic helper functions. Then comes the code for compiling a regular expression, and for showing the result of the compilation. The building phase converts a compiled regex to NFA states, and the automaton is run by the code in the following section. The only remaining brick is parsing the replacement text and performing the replacement. We are then ready for all the user functions. Finally, messages, and a little bit of tracing code.

41.2 Helpers

`__regex_int_eval:w` Access the primitive: performance is key here, so we do not use the slower route *via* `\int_eval:n`.

```
24105 \cs_new_eq:NN \__regex_int_eval:w \tex_numexpr:D
```

(End definition for `__regex_int_eval:w`.)

`__regex_standard_escapechar:` Make the `\escapechar` into the standard backslash.

```
24106 \cs_new_protected:Npn \__regex_standard_escapechar:
24107 { \int_set:Nn \tex_escapechar:D { '\ } }
```

(End definition for `__regex_standard_escapechar:.`)

`__regex_toks_use:w` Unpack a `\toks` given its number.

```
24108 \cs_new:Npn \__regex_toks_use:w { \tex_the:D \tex_toks:D }
```

(End definition for `__regex_toks_use:w`.)

`__regex_toks_clear:N` Empty a `\toks` or set it to a value, given its number.

```
\__regex_toks_set:Nn
\__regex_toks_set:No
24109 \cs_new_protected:Npn \__regex_toks_clear:N #1
24110 { \__regex_toks_set:Nn #1 { } }
24111 \cs_new_eq:NN \__regex_toks_set:Nn \tex_toks:D
24112 \cs_new_protected:Npn \__regex_toks_set:No #1
24113 { \tex_toks:D #1 \exp_after:wN }
```

(End definition for `__regex_toks_clear:N` and `__regex_toks_set:Nn`.)

`__regex_toks_memcpy:NNn` Copy #3 `\toks` registers from #2 onwards to #1 onwards, like C's `memcpy`.

```
24114 \cs_new_protected:Npn \__regex_toks_memcpy:NNn #1#2#3
24115 {
24116   \prg_replicate:nn {#3}
24117   {
24118     \tex_toks:D #1 = \tex_toks:D #2
24119     \int_incr:N #1
24120     \int_incr:N #2
24121   }
24122 }
```

(End definition for `__regex_toks_memcpy:NNn`.)

`__regex_toks_put_left:Nx` During the building phase we wish to add x-expanded material to `\toks`, either to the left or to the right. The expansion is done “by hand” for optimization (these operations are used quite a lot). The `Nn` version of `__regex_toks_put_right:Nx` is provided because it is more efficient than x-expanding with `\exp_not:n`.

```
24123 \cs_new_protected:Npn \__regex_toks_put_left:Nx #1#2
24124 {
24125   \cs_set_nopar:Npx \__regex_tmp:w { #2 }
24126   \tex_toks:D #1 \exp_after:wN \exp_after:wN \exp_after:wN
24127   { \exp_after:wN \__regex_tmp:w \tex_the:D \tex_toks:D #1 }
24128 }
24129 \cs_new_protected:Npn \__regex_toks_put_right:Nx #1#2
24130 {
24131   \cs_set_nopar:Npx \__regex_tmp:w {#2}
24132   \tex_toks:D #1 \exp_after:wN
```

```

24133     { \tex_the:D \tex_toks:D \exp_after:wN #1 \__regex_tmp:w }
24134   }
24135   \cs_new_protected:Npn \__regex_toks_put_right:Nn #1#2
24136   { \tex_toks:D #1 \exp_after:wN { \tex_the:D \tex_toks:D #1 #2 } }

```

(End definition for __regex_toks_put_left:Nx and __regex_toks_put_right:Nx.)

__regex_curr_cs_to_str: Expands to the string representation of the token (known to be a control sequence) at the current position \l__regex_curr_pos_int. It should only be used in x-expansion to avoid losing a leading space.

```

24137 \cs_new:Npn \__regex_curr_cs_to_str:
24138 {
24139   \exp_after:wN \exp_after:wN \exp_after:wN \cs_to_str:N
24140   \l__regex_curr_token_tl
24141 }

```

(End definition for __regex_curr_cs_to_str:.)

__regex_intarray_item:NnF Item of intarray, with a default value.

```

\__regex_intarray_item_aux:nNF
24142 \cs_new:Npn \__regex_intarray_item:NnF #1#2
24143 { \exp_args:Nf \__regex_intarray_item_aux:nNF { \int_eval:n {#2} } #1 }
24144 \cs_new:Npn \__regex_intarray_item_aux:nNF #1#2
24145 {
24146   \if_int_compare:w #1 > \c_zero_int
24147     \exp_after:wN \use_i:nn
24148   \else:
24149     \exp_after:wN \use_ii:nn
24150   \fi:
24151   { \__kernel_intarray_item:Nn #2 {#1} }
24152 }

```

(End definition for __regex_intarray_item:NnF and __regex_intarray_item_aux:nNF.)

__regex_maplike_break: Analogous to \tl_map_break:, this correctly exits \tl_map_inline:nn and similar constructions and jumps to the matching \prg_break_point:Nn __regex_maplike_break: { }.

```

24153 \cs_new:Npn \__regex_maplike_break:
24154 { \prg_map_break:Nn \__regex_maplike_break: { } }

```

(End definition for __regex_maplike_break:.)

41.2.1 Constants and variables

__regex_tmp:w Temporary function used for various short-term purposes.

```

24155 \cs_new:Npn \__regex_tmp:w { }

```

(End definition for __regex_tmp:w.)

\l__regex_internal_a_tl Temporary variables used for various purposes.

```

\l__regex_internal_b_tl
\l__regex_internal_a_int
\l__regex_internal_b_int
\l__regex_internal_c_int
\l__regex_internal_bool
\l__regex_internal_seq
\g__regex_internal_tl
24156 \tl_new:N \l__regex_internal_a_tl
24157 \tl_new:N \l__regex_internal_b_tl
24158 \int_new:N \l__regex_internal_a_int
24159 \int_new:N \l__regex_internal_b_int
24160 \int_new:N \l__regex_internal_c_int
24161 \bool_new:N \l__regex_internal_bool
24162 \seq_new:N \l__regex_internal_seq
24163 \tl_new:N \g__regex_internal_tl

```


(End definition for \l__regex_internal_a_t1 and others.)

`\l__regex_build_t1` This temporary variable is specifically for use with the `t1_build` machinery.

```
24164 \tl_new:N \l__regex_build_t1
```

(End definition for \l__regex_build_t1.)

`\c__regex_no_match_regex` This regular expression matches nothing, but is still a valid regular expression. We could use a failing assertion, but I went for an empty class. It is used as the initial value for regular expressions declared using `\regex_new:N`.

```
24165 \tl_const:Nn \c__regex_no_match_regex
24166 {
24167   \__regex_branch:n
24168   { \__regex_class:NnnnN \c_true_bool { } { 1 } { 0 } \c_true_bool }
24169 }
```

(End definition for \c__regex_no_match_regex.)

`\l__regex_balance_int` During this phase, `\l__regex_balance_int` counts the balance of begin-group and end-group character tokens which appear before a given point in the token list. This variable is also used to keep track of the balance in the replacement text.

```
24170 \int_new:N \l__regex_balance_int
```

(End definition for \l__regex_balance_int.)

41.2.2 Testing characters

```
\c__regex_ascii_min_int
\c__regex_ascii_max_control_int 24171 \int_const:Nn \c__regex_ascii_min_int { 0 }
\c__regex_ascii_max_int          24172 \int_const:Nn \c__regex_ascii_max_control_int { 31 }
24173 \int_const:Nn \c__regex_ascii_max_int { 127 }
```

(End definition for \c__regex_ascii_min_int, \c__regex_ascii_max_control_int, and \c__regex_ascii_max_int.)

```
\c__regex_ascii_lower_int
24174 \int_const:Nn \c__regex_ascii_lower_int { 'a - 'A }
```

(End definition for \c__regex_ascii_lower_int.)

41.2.3 Internal auxiliaries

`\q__regex_recursion_stop` Internal recursion quarks.

```
24175 \quark_new:N \q__regex_recursion_stop
```

(End definition for \q__regex_recursion_stop.)

`__regex_use_none_delimit_by_q_recursion_stop:w` Functions to gobble up to a quark.

```
\__regex_use_i_delimit_by_q_recursion_stop:nw 24176 \cs_new:Npn \__regex_use_none_delimit_by_q_recursion_stop:w
24177   #1 \q__regex_recursion_stop { }
24178 \cs_new:Npn \__regex_use_i_delimit_by_q_recursion_stop:nw
24179   #1 #2 \q__regex_recursion_stop {#1}
```

(End definition for __regex_use_none_delimit_by_q_recursion_stop:w and __regex_use_i_delimit_by_q_recursion_stop:nw.)

`\q__regex_nil` Internal quarks.

```
24180 \quark_new:N \q__regex_nil
```

(End definition for `\q__regex_nil`.)

`__regex_quark_if_nil_p:n` Branching quark conditional.

```
\__regex_quark_if_nil:nTF 24181 \__kernel_quark_new_conditional:Nn \__regex_quark_if_nil:N { F }
```

(End definition for `__regex_quark_if_nil:nTF`.)

`__regex_break_point:TF` When testing whether a character of the query token list matches a given character class in the regular expression, we often have to test it against several ranges of characters, checking if any one of those matches. This is done with a structure like

```
\__regex_break_point:TF {<true code>} {<false code>}
```

If any of the tests succeeds, it calls `__regex_break_true:w`, which cleans up and leaves `<true code>` in the input stream. Otherwise, `__regex_break_point:TF` leaves the `<false code>` in the input stream.

```
24182 \cs_new_protected:Npn \__regex_break_true:w
24183   #1 \__regex_break_point:TF #2 #3 {#2}
24184 \cs_new_protected:Npn \__regex_break_point:TF #1 #2 { #2 }
```

(End definition for `__regex_break_point:TF` and `__regex_break_true:w`.)

`__regex_item_reverse:n` This function makes showing regular expressions easier, and lets us define `\D` in terms of `\d` for instance. There is a subtlety: the end of the query is marked by `-2`, and thus matches `\D` and other negated properties; this case is caught by another part of the code.

```
24185 \cs_new_protected:Npn \__regex_item_reverse:n #1
24186   {
24187     #1
24188     \__regex_break_point:TF { } \__regex_break_true:w
24189   }
```

(End definition for `__regex_item_reverse:n`.)

`__regex_item_caseful_equal:n` Simple comparisons triggering `__regex_break_true:w` when true.

```
\__regex_item_caseful_range:nn 24190 \cs_new_protected:Npn \__regex_item_caseful_equal:n #1
24191   {
24192     \if_int_compare:w #1 = \l__regex_curr_char_int
24193       \exp_after:wN \__regex_break_true:w
24194     \fi:
24195   }
24196 \cs_new_protected:Npn \__regex_item_caseful_range:nn #1 #2
24197   {
24198     \reverse_if:N \if_int_compare:w #1 > \l__regex_curr_char_int
24199     \reverse_if:N \if_int_compare:w #2 < \l__regex_curr_char_int
24200     \exp_after:wN \exp_after:wN \exp_after:wN \__regex_break_true:w
24201     \fi:
24202     \fi:
24203   }
```

(End definition for `__regex_item_caseful_equal:n` and `__regex_item_caseful_range:nn`.)

_regex_item_caseless_equal:n
_regex_item_caseless_range:nn

For caseless matching, we perform the test both on the `curr_char` and on the `case_changed_char`. Before doing the second set of tests, we make sure that `case_changed_char` has been computed.

```

24204 \cs_new_protected:Npn \__regex_item_caseless_equal:n #1
24205 {
24206   \if_int_compare:w #1 = \l__regex_curr_char_int
24207   \exp_after:wN \__regex_break_true:w
24208   \fi:
24209   \if_int_compare:w \l__regex_case_changed_char_int = \c_max_int
24210   \__regex_compute_case_changed_char:
24211   \fi:
24212   \if_int_compare:w #1 = \l__regex_case_changed_char_int
24213   \exp_after:wN \__regex_break_true:w
24214   \fi:
24215 }
24216 \cs_new_protected:Npn \__regex_item_caseless_range:nn #1 #2
24217 {
24218   \reverse_if:N \if_int_compare:w #1 > \l__regex_curr_char_int
24219   \reverse_if:N \if_int_compare:w #2 < \l__regex_curr_char_int
24220   \exp_after:wN \exp_after:wN \exp_after:wN \__regex_break_true:w
24221   \fi:
24222   \fi:
24223   \if_int_compare:w \l__regex_case_changed_char_int = \c_max_int
24224   \__regex_compute_case_changed_char:
24225   \fi:
24226   \reverse_if:N \if_int_compare:w #1 > \l__regex_case_changed_char_int
24227   \reverse_if:N \if_int_compare:w #2 < \l__regex_case_changed_char_int
24228   \exp_after:wN \exp_after:wN \exp_after:wN \__regex_break_true:w
24229   \fi:
24230   \fi:
24231 }

```

(End definition for _regex_item_caseless_equal:n and _regex_item_caseless_range:nn.)

_regex_compute_case_changed_char:

This function is called when `\l__regex_case_changed_char_int` has not yet been computed (or rather, when it is set to the marker value `\c_max_int`). If the current character code is in the range [65, 90] (upper-case), then add 32, making it lowercase. If it is in the lower-case letter range [97, 122], subtract 32.

```

24232 \cs_new_protected:Npn \__regex_compute_case_changed_char:
24233 {
24234   \int_set_eq:NN \l__regex_case_changed_char_int \l__regex_curr_char_int
24235   \if_int_compare:w \l__regex_curr_char_int > 'Z \exp_stop_f:
24236   \if_int_compare:w \l__regex_curr_char_int > 'z \exp_stop_f: \else:
24237   \if_int_compare:w \l__regex_curr_char_int < 'a \exp_stop_f: \else:
24238   \int_sub:Nn \l__regex_case_changed_char_int
24239   { \c__regex_ascii_lower_int }
24240   \fi:
24241   \fi:
24242   \else:
24243   \if_int_compare:w \l__regex_curr_char_int < 'A \exp_stop_f: \else:
24244   \int_add:Nn \l__regex_case_changed_char_int
24245   { \c__regex_ascii_lower_int }
24246   \fi:
24247   \fi:

```

24248 }

(End definition for `_regex_compute_case_changed_char:.`)

`_regex_item_equal:n` Those must always be defined to expand to a `caseful` (default) or `caseless` version, and not be protected: they must expand when compiling, to hard-code which tests are caseless or caseful.

24249 `\cs_new_eq:NN _regex_item_equal:n ?`

24250 `\cs_new_eq:NN _regex_item_range:nn ?`

(End definition for `_regex_item_equal:n` and `_regex_item_range:nn`.)

`_regex_item_catcode:nT` The argument is a sum of powers of 4 with exponents given by the allowed category codes (between 0 and 13). Dividing by a given power of 4 gives an odd result if and only if that category code is allowed. If the catcode does not match, then skip the character code tests which follow.

24251 `\cs_new_protected:Npn _regex_item_catcode:`

24252 {

24253 "

24254 `\if_case:w \l__regex_curr_catcode_int`

24255 `1 \or: 4 \or: 10 \or: 40`

24256 `\or: 100 \or: \or: 1000 \or: 4000`

24257 `\or: 10000 \or: \or: 100000 \or: 400000`

24258 `\or: 1000000 \or: 4000000 \else: 1*0`

24259 `\fi:`

24260 }

24261 `\cs_new_protected:Npn _regex_item_catcode:nT #1`

24262 {

24263 `\if_int_odd:w \int_eval:n { #1 / _regex_item_catcode: } \exp_stop_f:`

24264 `\exp_after:wN \use:n`

24265 `\else:`

24266 `\exp_after:wN \use_none:n`

24267 `\fi:`

24268 }

24269 `\cs_new_protected:Npn _regex_item_catcode_reverse:nT #1#2`

24270 `{ _regex_item_catcode:nT {#1} { _regex_item_reverse:n {#2} } }`

(End definition for `_regex_item_catcode:nT`, `_regex_item_catcode_reverse:nT`, and `_regex_item_catcode:.`)

`_regex_item_exact:nn` This matches an exact `<category>-<character code>` pair, or an exact control sequence, more precisely one of several possible control sequences, separated by `\scan_stop:.`

`_regex_item_exact_cs:n`

24271 `\cs_new_protected:Npn _regex_item_exact:nn #1#2`

24272 {

24273 `\if_int_compare:w #1 = \l__regex_curr_catcode_int`

24274 `\if_int_compare:w #2 = \l__regex_curr_char_int`

24275 `\exp_after:wN \exp_after:wN \exp_after:wN _regex_break_true:w`

24276 `\fi:`

24277 `\fi:`

24278 }

24279 `\cs_new_protected:Npn _regex_item_exact_cs:n #1`

24280 {

24281 `\int_compare:nNnTF \l__regex_curr_catcode_int = 0`

24282 {

```

24283     \_kernel_tl_set:Nx \l__regex_internal_a_tl
24284     { \scan_stop: \__regex_curr_cs_to_str: \scan_stop: }
24285     \tl_if_in:noTF { \scan_stop: #1 \scan_stop: }
24286     \l__regex_internal_a_tl
24287     { \__regex_break_true:w } { }
24288   }
24289   { }
24290 }

```

(End definition for `__regex_item_exact:nn` and `__regex_item_exact_cs:n`.)

`__regex_item_cs:n` Match a control sequence (the argument is a compiled regex). First test the catcode of the current token to be zero. Then perform the matching test, and break if the csname indeed matches.

```

24291 \cs_new_protected:Npn \__regex_item_cs:n #1
24292 {
24293   \int_compare:nNnT \l__regex_curr_catcode_int = 0
24294   {
24295     \group_begin:
24296     \__regex_single_match:
24297     \__regex_disable_submatches:
24298     \__regex_build_for_cs:n {#1}
24299     \bool_set_eq:NN \l__regex_saved_success_bool
24300     \g__regex_success_bool
24301     \exp_args:Nx \__regex_match_cs:n { \__regex_curr_cs_to_str: }
24302     \if_meaning:w \c_true_bool \g__regex_success_bool
24303     \group_insert_after:N \__regex_break_true:w
24304     \fi:
24305     \bool_gset_eq:NN \g__regex_success_bool
24306     \l__regex_saved_success_bool
24307   \group_end:
24308   }
24309 }

```

(End definition for `__regex_item_cs:n`.)

41.2.4 Character property tests

`__regex_prop_d:` Character property tests for `\d`, `\W`, *etc.* These character properties are not affected by the `(?i)` option. The characters recognized by each one are as follows: `\d`=[0-9], `\w`=[0-9A-Z_a-z], `\s`=[\ \^I\^J\^L\^M], `\h`=[\ \^I], `\v`=[\^J\^M], and the upper case counterparts match anything that the lower case does not match. The order in which the various tests appear is optimized for usual mostly lower case letter text.

```

24310 \cs_new_protected:Npn \__regex_prop_d:
24311 { \__regex_item_caseful_range:nn { '0 } { '9 } }
24312 \cs_new_protected:Npn \__regex_prop_h:
24313 {
24314   \__regex_item_caseful_equal:n { '\ }
24315   \__regex_item_caseful_equal:n { '\^I }
24316 }
24317 \cs_new_protected:Npn \__regex_prop_s:
24318 {
24319   \__regex_item_caseful_equal:n { '\ }
24320   \__regex_item_caseful_equal:n { '\^I }

```

```

24321     \_regex_item_caseful_equal:n { '\^J }
24322     \_regex_item_caseful_equal:n { '\^L }
24323     \_regex_item_caseful_equal:n { '\^M }
24324 }
24325 \cs_new_protected:Npn \_regex_prop_v:
24326 { \_regex_item_caseful_range:nn { '\^J } { '\^M } } % lf, vtab, ff, cr
24327 \cs_new_protected:Npn \_regex_prop_w:
24328 {
24329     \_regex_item_caseful_range:nn { 'a } { 'z }
24330     \_regex_item_caseful_range:nn { 'A } { 'Z }
24331     \_regex_item_caseful_range:nn { '0 } { '9 }
24332     \_regex_item_caseful_equal:n { '_' }
24333 }
24334 \cs_new_protected:Npn \_regex_prop_N:
24335 {
24336     \_regex_item_reverse:n
24337     { \_regex_item_caseful_equal:n { '\^J } }
24338 }

```

(End definition for _regex_prop_d: and others.)

```

\_regex_posix_alnum: POSIX properties. No surprise.
\_regex_posix_alpha:
\_regex_posix_ascii:
\_regex_posix_blank:
\_regex_posix_cntrl:
\_regex_posix_digit:
\_regex_posix_graph:
\_regex_posix_lower:
\_regex_posix_print:
\_regex_posix_punct:
\_regex_posix_space:
\_regex_posix_upper:
\_regex_posix_word:
\_regex_posix_xdigit:
24339 \cs_new_protected:Npn \_regex_posix_alnum:
24340 { \_regex_posix_alpha: \_regex_posix_digit: }
24341 \cs_new_protected:Npn \_regex_posix_alpha:
24342 { \_regex_posix_lower: \_regex_posix_upper: }
24343 \cs_new_protected:Npn \_regex_posix_ascii:
24344 {
24345     \_regex_item_caseful_range:nn
24346     \c__regex_ascii_min_int
24347     \c__regex_ascii_max_int
24348 }
24349 \cs_new_eq:NN \_regex_posix_blank: \_regex_prop_h:
24350 \cs_new_protected:Npn \_regex_posix_cntrl:
24351 {
24352     \_regex_item_caseful_range:nn
24353     \c__regex_ascii_min_int
24354     \c__regex_ascii_max_control_int
24355     \_regex_item_caseful_equal:n \c__regex_ascii_max_int
24356 }
24357 \cs_new_eq:NN \_regex_posix_digit: \_regex_prop_d:
24358 \cs_new_protected:Npn \_regex_posix_graph:
24359 { \_regex_item_caseful_range:nn { '!' } { '~ } }
24360 \cs_new_protected:Npn \_regex_posix_lower:
24361 { \_regex_item_caseful_range:nn { 'a' } { 'z' } }
24362 \cs_new_protected:Npn \_regex_posix_print:
24363 { \_regex_item_caseful_range:nn { '\' } { '~ } }
24364 \cs_new_protected:Npn \_regex_posix_punct:
24365 {
24366     \_regex_item_caseful_range:nn { '!' } { '/' }
24367     \_regex_item_caseful_range:nn { ':' } { '@' }
24368     \_regex_item_caseful_range:nn { '[' } { '[' }
24369     \_regex_item_caseful_range:nn { '{' } { '~ }
24370 }

```

```

24371 \cs_new_protected:Npn \__regex_posix_space:
24372 {
24373   \__regex_item_caseful_equal:n { '\ }
24374   \__regex_item_caseful_range:nn { '\^I } { '\^M }
24375 }
24376 \cs_new_protected:Npn \__regex_posix_upper:
24377 { \__regex_item_caseful_range:nn { 'A } { 'Z } }
24378 \cs_new_eq:NN \__regex_posix_word: \__regex_prop_w:
24379 \cs_new_protected:Npn \__regex_posix_xdigit:
24380 {
24381   \__regex_posix_digit:
24382   \__regex_item_caseful_range:nn { 'A } { 'F }
24383   \__regex_item_caseful_range:nn { 'a } { 'f }
24384 }

```

(End definition for `__regex_posix_alnum:` and others.)

41.2.5 Simple character escape

Before actually parsing the regular expression or the replacement text, we go through them once, converting `\n` to the character 10, *etc.* In this pass, we also convert any special character (`*`, `?`, `{`, *etc.*) or escaped alphanumeric character into a marker indicating that this was a special sequence, and replace escaped special characters and non-escaped alphanumeric characters by markers indicating that those were “raw” characters. The rest of the code can then avoid caring about escaping issues (those can become quite complex to handle in combination with ranges in character classes).

Usage: `__regex_escape_use:nnnn` *<inline 1>* *<inline 2>* *<inline 3>* *<{token list}>*
The *<token list>* is converted to a string, then read from left to right, interpreting backslashes as escaping the next character. Unescaped characters are fed to the function *<inline 1>*, and escaped characters are fed to the function *<inline 2>* within an *x*-expansion context (typically those functions perform some tests on their argument to decide how to output them). The escape sequences `\a`, `\e`, `\f`, `\n`, `\r`, `\t` and `\x` are recognized, and those are replaced by the corresponding character, then fed to *<inline 3>*. The result is then left in the input stream. Spaces are ignored unless escaped.

The conversion is done within an *x*-expanding assignment.

`__regex_escape_use:nnnn` The result is built in `\l__regex_internal_a_tl`, which is then left in the input stream. Tracing code is added as appropriate inside this token list. Go through #4 once, applying #1, #2, or #3 as relevant to each character (after de-escaping it).

```

24385 \cs_new_protected:Npn \__regex_escape_use:nnnn #1#2#3#4
24386 {
24387   \group_begin:
24388   \tl_clear:N \l__regex_internal_a_tl
24389   \cs_set:Npn \__regex_escape_unescaped:N ##1 { #1 }
24390   \cs_set:Npn \__regex_escape_escaped:N ##1 { #2 }
24391   \cs_set:Npn \__regex_escape_raw:N ##1 { #3 }
24392   \__regex_standard_escapechar:
24393   \__kernel_tl_gset:Nx \g__regex_internal_tl
24394   { \__kernel_str_to_other_fast:n {#4} }
24395   \tl_put_right:Nx \l__regex_internal_a_tl
24396   {
24397     \exp_after:wN \__regex_escape_loop:N \g__regex_internal_tl
24398     { break } \prg_break_point:

```

```

24399     }
24400     \exp_after:wN
24401     \group_end:
24402     \l__regex_internal_a_tl
24403 }

```

(End definition for __regex_escape_use:nnnn.)

__regex_escape_loop:N __regex_escape_loop:N reads one character: if it is special (space, backslash, or end-marker), perform the associated action, otherwise it is simply an unescaped character. After a backslash, the same is done, but unknown characters are “escaped”.

```

24404 \cs_new:Npn \__regex_escape_loop:N #1
24405 {
24406     \cs_if_exist_use:cF { __regex_escape\_token_to_str:N #1:w }
24407     { \__regex_escape_unescaped:N #1 }
24408     \__regex_escape_loop:N
24409 }
24410 \cs_new:cpn { __regex_escape\_c_backslash_str :w }
24411     \__regex_escape_loop:N #1
24412 {
24413     \cs_if_exist_use:cF { __regex_escape\_token_to_str:N #1:w }
24414     { \__regex_escape_escaped:N #1 }
24415     \__regex_escape_loop:N
24416 }

```

(End definition for __regex_escape_loop:N and __regex_escape_w.)

__regex_escape_unescaped:N Those functions are never called before being given a new meaning, so their definitions here don’t matter.

```

\__regex_escape_escaped:N
\__regex_escape_raw:N
24417 \cs_new_eq:NN \__regex_escape_unescaped:N ?
24418 \cs_new_eq:NN \__regex_escape_escaped:N ?
24419 \cs_new_eq:NN \__regex_escape_raw:N ?

```

(End definition for __regex_escape_unescaped:N, __regex_escape_escaped:N, and __regex_escape_raw:N.)

__regex_escape_break:w The loop is ended upon seeing the end-marker “break”, with an error if the string ended in a backslash. Spaces are ignored, and \a, \e, \f, \n, \r, \t take their meaning here.

```

\__regex_escape_/break:w
\__regex_escape_/a:w
\__regex_escape_/e:w
\__regex_escape_/f:w
\__regex_escape_/n:w
\__regex_escape_/r:w
\__regex_escape_/t:w
\__regex_escape_\_w
24420 \cs_new_eq:NN \__regex_escape_break:w \prg_break:
24421 \cs_new:cpn { __regex_escape_/break:w }
24422 {
24423     \__kernel_msg_expandable_error:nn { kernel } { trailing-backslash }
24424     \prg_break:
24425 }
24426 \cs_new:cpn { __regex_escape\_~:w } { }
24427 \cs_new:cpx { __regex_escape_/a:w }
24428 { \exp_not:N \__regex_escape_raw:N \iow_char:N \^^G }
24429 \cs_new:cpx { __regex_escape_/t:w }
24430 { \exp_not:N \__regex_escape_raw:N \iow_char:N \^^I }
24431 \cs_new:cpx { __regex_escape_/n:w }
24432 { \exp_not:N \__regex_escape_raw:N \iow_char:N \^^J }
24433 \cs_new:cpx { __regex_escape_/f:w }
24434 { \exp_not:N \__regex_escape_raw:N \iow_char:N \^^L }
24435 \cs_new:cpx { __regex_escape_/r:w }

```



```

24436 { \exp_not:N \__regex_escape_raw:N \iow_char:N \^^M }
24437 \cs_new:cpx { __regex_escape_/e:w }
24438 { \exp_not:N \__regex_escape_raw:N \iow_char:N \^^[ }

```

(End definition for __regex_escape_break:w and others.)

```

\__regex_escape_/x:w
\__regex_escape_x_end:w
\__regex_escape_x_large:n

```

When \x is encountered, __regex_escape_x_test:N is responsible for grabbing some hexadecimal digits, and feeding the result to __regex_escape_x_end:w. If the number is too big interrupt the assignment and produce an error, otherwise call __regex_escape_raw:N on the corresponding character token.

```

24439 \cs_new:cpn { __regex_escape_/x:w } \__regex_escape_loop:N
24440 {
24441   \exp_after:wN \__regex_escape_x_end:w
24442   \int_value:w "0 \__regex_escape_x_test:N
24443 }
24444 \cs_new:Npn \__regex_escape_x_end:w #1 ;
24445 {
24446   \int_compare:nNnTF {#1} > \c_max_char_int
24447   {
24448     \__kernel_msg_expandable_error:nnff { kernel } { x-overflow }
24449     {#1} { \int_to_Hex:n {#1} }
24450   }
24451   {
24452     \exp_last_unbraced:Nf \__regex_escape_raw:N
24453     { \char_generate:nn {#1} { 12 } }
24454   }
24455 }

```

(End definition for __regex_escape_/x:w, __regex_escape_x_end:w, and __regex_escape_x_large:n.)

```

\__regex_escape_x_test:N
\__regex_escape_x_testii:N

```

Find out whether the first character is a left brace (allowing any number of hexadecimal digits), or not (allowing up to two hexadecimal digits). We need to check for the end-of-string marker. Eventually, call either __regex_escape_x_loop:N or __regex_escape_x:N.

```

24456 \cs_new:Npn \__regex_escape_x_test:N #1
24457 {
24458   \str_if_eq:nnTF {#1} { break } { ; }
24459   {
24460     \if_charcode:w \c_space_token #1
24461     \exp_after:wN \__regex_escape_x_test:N
24462   \else:
24463     \exp_after:wN \__regex_escape_x_testii:N
24464     \exp_after:wN #1
24465   \fi:
24466   }
24467 }
24468 \cs_new:Npn \__regex_escape_x_testii:N #1
24469 {
24470   \if_charcode:w \c_left_brace_str #1
24471   \exp_after:wN \__regex_escape_x_loop:N
24472 \else:
24473   \__regex_hexadecimal_use:NTF #1
24474   { \exp_after:wN \__regex_escape_x:N }
24475   { ; \exp_after:wN \__regex_escape_loop:N \exp_after:wN #1 }

```

```

24476     \fi:
24477   }

```

(End definition for `_regex_escape_x_test:N` and `_regex_escape_x_testii:N`.)

`_regex_escape_x:N` This looks for the second digit in the unbraced case.

```

24478 \cs_new:Npn \_regex_escape_x:N #1
24479 {
24480   \str_if_eq:nnTF {#1} { break } { ; }
24481   {
24482     \_regex_hexadecimal_use:NTF #1
24483     { ; \_regex_escape_loop:N }
24484     { ; \_regex_escape_loop:N #1 }
24485   }
24486 }

```

(End definition for `_regex_escape_x:N`.)

`_regex_escape_x_loop:N` Grab hexadecimal digits, skip spaces, and at the end, check that there is a right brace,
`_regex_escape_x_loop_error:` otherwise raise an error outside the assignment.

```

24487 \cs_new:Npn \_regex_escape_x_loop:N #1
24488 {
24489   \str_if_eq:nnTF {#1} { break }
24490   { ; \_regex_escape_x_loop_error:n { } {#1} }
24491   {
24492     \_regex_hexadecimal_use:NTF #1
24493     { \_regex_escape_x_loop:N }
24494     {
24495       \token_if_eq_charcode:NNTF \c_space_token #1
24496       { \_regex_escape_x_loop:N }
24497       {
24498         ;
24499         \exp_after:wN
24500         \token_if_eq_charcode:NNTF \c_right_brace_str #1
24501         { \_regex_escape_loop:N }
24502         { \_regex_escape_x_loop_error:n {#1} }
24503       }
24504     }
24505   }
24506 }
24507 \cs_new:Npn \_regex_escape_x_loop_error:n #1
24508 {
24509   \__kernel_msg_expandable_error:nnn { kernel } { x-missing-rbrace } {#1}
24510   \_regex_escape_loop:N #1
24511 }

```

(End definition for `_regex_escape_x_loop:N` and `_regex_escape_x_loop_error:.`)

`_regex_hexadecimal_use:NTF` T_EX detects uppercase hexadecimal digits for us but not the lowercase letters, which we need to detect and replace by their uppercase counterpart.

```

24512 \prg_new_conditional:Npnn \_regex_hexadecimal_use:N #1 { TF }
24513 {
24514   \if_int_compare:w 1 < "1 \token_to_str:N #1 \exp_stop_f:
24515   #1 \prg_return_true:
24516   \else:

```

```

24517     \if_case:w
24518         \int_eval:n { \exp_after:wN ‘ \token_to_str:N #1 - ‘a }
24519         A
24520     \or: B
24521     \or: C
24522     \or: D
24523     \or: E
24524     \or: F
24525     \else:
24526         \prg_return_false:
24527         \exp_after:wN \use_none:n
24528     \fi:
24529     \prg_return_true:
24530 \fi:
24531 }

```

(End definition for `__regex_hexadecimal_use:NTF`.)

`__regex_char_if_alphanumeric:NTF`
`__regex_char_if_special:NTF`

These two tests are used in the first pass when parsing a regular expression. That pass is responsible for finding escaped and non-escaped characters, and recognizing which ones have special meanings and which should be interpreted as “raw” characters. Namely,

- alphanumeric are “raw” if they are not escaped, and may have a special meaning when escaped;
- non-alphanumeric printable ascii characters are “raw” if they are escaped, and may have a special meaning when not escaped;
- characters other than printable ascii are always “raw”.

The code is ugly, and highly based on magic numbers and the ascii codes of characters. This is mostly unavoidable for performance reasons. Maybe the tests can be optimized a little bit more. Here, “alphanumeric” means 0–9, A–Z, a–z; “special” character means non-alphanumeric but printable ascii, from space (hex 20) to del (hex 7E).

```

24532 \prg_new_conditional:Npnn \__regex_char_if_special:N #1 { TF }
24533 {
24534     \if_int_compare:w ‘#1 > ‘Z \exp_stop_f:
24535     \if_int_compare:w ‘#1 > ‘z \exp_stop_f:
24536     \if_int_compare:w ‘#1 < \c__regex_ascii_max_int
24537     \prg_return_true: \else: \prg_return_false: \fi:
24538     \else:
24539     \if_int_compare:w ‘#1 < ‘a \exp_stop_f:
24540     \prg_return_true: \else: \prg_return_false: \fi:
24541     \fi:
24542     \else:
24543     \if_int_compare:w ‘#1 > ‘9 \exp_stop_f:
24544     \if_int_compare:w ‘#1 < ‘A \exp_stop_f:
24545     \prg_return_true: \else: \prg_return_false: \fi:
24546     \else:
24547     \if_int_compare:w ‘#1 < ‘0 \exp_stop_f:
24548     \if_int_compare:w ‘#1 < ‘\ \exp_stop_f:
24549     \prg_return_false: \else: \prg_return_true: \fi:
24550     \else: \prg_return_false: \fi:
24551     \fi:
24552 \fi:

```

```

24553     }
24554 \prg_new_conditional:Npnn \__regex_char_if_alphanumeric:N #1 { TF }
24555 {
24556     \if_int_compare:w '#1 > 'Z \exp_stop_f:
24557     \if_int_compare:w '#1 > 'z \exp_stop_f:
24558     \prg_return_false:
24559     \else:
24560     \if_int_compare:w '#1 < 'a \exp_stop_f:
24561     \prg_return_false: \else: \prg_return_true: \fi:
24562     \fi:
24563     \else:
24564     \if_int_compare:w '#1 > '9 \exp_stop_f:
24565     \if_int_compare:w '#1 < 'A \exp_stop_f:
24566     \prg_return_false: \else: \prg_return_true: \fi:
24567     \else:
24568     \if_int_compare:w '#1 < '0 \exp_stop_f:
24569     \prg_return_false: \else: \prg_return_true: \fi:
24570     \fi:
24571     \fi:
24572 }

```

(End definition for __regex_char_if_alphanumeric:N and __regex_char_if_special:N.)

41.3 Compiling

A regular expression starts its life as a string of characters. In this section, we convert it to internal instructions, resulting in a “compiled” regular expression. This compiled expression is then turned into states of an automaton in the building phase. Compiled regular expressions consist of the following:

- __regex_class:NnnN <boolean> {<tests>} {<min>} {<more>} <lazyness>
- __regex_group:nnnN {<branches>} {<min>} {<more>} <lazyness>, also __regex_group_no_capture:nnnN and __regex_group_resetting:nnnN with the same syntax.
- __regex_branch:n {<contents>}
- __regex_command_K:
- __regex_assertion:Nn <boolean> {<assertion test>}, where the <assertion test> is __regex_b_test: or __regex_Z_test: or __regex_A_test: or __regex_G_test:

Tests can be the following:

- __regex_item_caseful_equal:n {<char code>}
- __regex_item_caseless_equal:n {<char code>}
- __regex_item_caseful_range:nn {<min>} {<max>}
- __regex_item_caseless_range:nn {<min>} {<max>}
- __regex_item_catcode:nT {<catcode bitmap>} {<tests>}

- `__regex_item_catcode_reverse:nT {⟨catcode bitmap⟩} {⟨tests⟩}`
- `__regex_item_reverse:n {⟨tests⟩}`
- `__regex_item_exact:nn {⟨catcode⟩} {⟨char code⟩}`
- `__regex_item_exact_cs:n {⟨csnames⟩}`, more precisely given as `⟨csname⟩ \scan_stop: ⟨csname⟩ \scan_stop: ⟨csname⟩` and so on in a brace group.
- `__regex_item_cs:n {⟨compiled regex⟩}`

41.3.1 Variables used when compiling

`\l__regex_group_level_int` We make sure to open the same number of groups as we close.

```
24573 \int_new:N \l__regex_group_level_int
```

(End definition for `\l__regex_group_level_int`.)

`\l__regex_mode_int` While compiling, ten modes are recognized, labelled -63 , -23 , -6 , -2 , 0 , 2 , 3 , 6 , 23 , 63 .
`\c__regex_cs_in_class_mode_int` See section 41.3.3. We only define some of these as constants.

```
24574 \int_new:N \l__regex_mode_int
24575 \int_const:Nn \c__regex_cs_in_class_mode_int { -6 }
24576 \int_const:Nn \c__regex_cs_mode_int { -2 }
24577 \int_const:Nn \c__regex_outer_mode_int { 0 }
24578 \int_const:Nn \c__regex_catcode_mode_int { 2 }
24579 \int_const:Nn \c__regex_class_mode_int { 3 }
24580 \int_const:Nn \c__regex_catcode_in_class_mode_int { 6 }
```

(End definition for `\l__regex_mode_int` and others.)

`\l__regex_catcodes_int` We wish to allow constructions such as `\c[~BE](. . \cL[a-z] . .)`, where the outer catcode test applies to the whole group, but is superseded by the inner catcode test. For this to work, we need to keep track of lists of allowed category codes: `\l__regex_catcodes_int` and `\l__regex_default_catcodes_int` are bitmaps, sums of 4^c , for all allowed catcodes c . The latter is local to each capturing group, and we reset `\l__regex_catcodes_int` to that value after each character or class, changing it only when encountering a `\c` escape. The boolean records whether the list of categories of a catcode test has to be inverted: compare `\c[~BE]` and `\c[BE]`.

```
24581 \int_new:N \l__regex_catcodes_int
24582 \int_new:N \l__regex_default_catcodes_int
24583 \bool_new:N \l__regex_catcodes_bool
```

(End definition for `\l__regex_catcodes_int`, `\l__regex_default_catcodes_int`, and `\l__regex_catcodes_bool`.)

`\c__regex_catcode_C_int` Constants: 4^c for each category, and the sum of all powers of 4.

```
24584 \int_const:Nn \c__regex_catcode_C_int { "1 }
24585 \int_const:Nn \c__regex_catcode_B_int { "4 }
24586 \int_const:Nn \c__regex_catcode_E_int { "10 }
24587 \int_const:Nn \c__regex_catcode_M_int { "40 }
24588 \int_const:Nn \c__regex_catcode_T_int { "100 }
24589 \int_const:Nn \c__regex_catcode_P_int { "1000 }
24590 \int_const:Nn \c__regex_catcode_U_int { "4000 }
24591 \int_const:Nn \c__regex_catcode_D_int { "10000 }
24592 \int_const:Nn \c__regex_catcode_S_int { "100000 }
```

`\c__regex_catcode_O_int`
`\c__regex_catcode_A_int`
`\c__regex_all_catcodes_int`

```

24593 \int_const:Nn \c__regex_catcode_L_int { "400000 }
24594 \int_const:Nn \c__regex_catcode_O_int { "1000000 }
24595 \int_const:Nn \c__regex_catcode_A_int { "4000000 }
24596 \int_const:Nn \c__regex_all_catcodes_int { "5515155 }

```

(End definition for \c__regex_catcode_C_int and others.)

\l__regex_internal_regex The compilation step stores its result in this variable.

```

24597 \cs_new_eq:NN \l__regex_internal_regex \c__regex_no_match_regex

```

(End definition for \l__regex_internal_regex.)

\l__regex_show_prefix_seq This sequence holds the prefix that makes up the line displayed to the user. The various items must be removed from the right, which is tricky with a token list, hence we use a sequence.

```

24598 \seq_new:N \l__regex_show_prefix_seq

```

(End definition for \l__regex_show_prefix_seq.)

\l__regex_show_lines_int A hack. To know whether a given class has a single item in it or not, we count the number of lines when showing the class.

```

24599 \int_new:N \l__regex_show_lines_int

```

(End definition for \l__regex_show_lines_int.)

41.3.2 Generic helpers used when compiling

__regex_two_if_eq:NNNNTF Used to compare pairs of things like __regex_compile_special:N ? together. It's often inconvenient to get the catcodes of the character to match so we just compare the character code. Besides, the expanding behaviour of \if:w is very useful as that means we can use \c_left_brace_str and the like.

```

24600 \prg_new_conditional:Npnn \__regex_two_if_eq:NNNN #1#2#3#4 { TF }
24601 {
24602   \if_meaning:w #1 #3
24603   \if:w #2 #4
24604     \prg_return_true:
24605   \else:
24606     \prg_return_false:
24607   \fi:
24608 \else:
24609   \prg_return_false:
24610 \fi:
24611 }

```

(End definition for __regex_two_if_eq:NNNNTF.)

__regex_get_digits:NTFw If followed by some raw digits, collect them one by one in the integer variable #1, and take the true branch. Otherwise, take the false branch.

```

24612 \cs_new_protected:Npn \__regex_get_digits:NTFw #1#2#3#4#5
24613 {
24614   \__regex_if_raw_digit:NNTF #4 #5
24615   { #1 = #5 \__regex_get_digits_loop:nw {#2} }
24616   { #3 #4 #5 }
24617 }

```

```

24618 \cs_new:Npn \__regex_get_digits_loop:nw #1#2#3
24619 {
24620   \__regex_if_raw_digit:NNTF #2 #3
24621   { #3 \__regex_get_digits_loop:nw {#1} }
24622   { \scan_stop: #1 #2 #3 }
24623 }

```

(End definition for __regex_get_digits:NNTFw and __regex_get_digits_loop:w.)

__regex_if_raw_digit:NNTF Test used when grabbing digits for the {m,n} quantifier. It only accepts non-escaped digits.

```

24624 \prg_new_conditional:Npnn \__regex_if_raw_digit:NN #1#2 { TF }
24625 {
24626   \if_meaning:w \__regex_compile_raw:N #1
24627   \if_int_compare:w 1 < 1 #2 \exp_stop_f:
24628   \prg_return_true:
24629   \else:
24630   \prg_return_false:
24631   \fi:
24632   \else:
24633   \prg_return_false:
24634   \fi:
24635 }

```

(End definition for __regex_if_raw_digit:NNTF.)

41.3.3 Mode

When compiling the NFA corresponding to a given regex string, we can be in ten distinct modes, which we label by some magic numbers:

- 6 [\c{...}] control sequence in a class,
- 2 \c{...} control sequence,
- 0 ... outer,
- 2 \c... catcode test,
- 6 [\c...] catcode test in a class,
- 63 [\c{[...]}] class inside mode -6,
- 23 \c{[...]} class inside mode -2,
- 3 [...] class inside mode 0,
- 23 \c[...] class inside mode 2,
- 63 [\c[...]] class inside mode 6.

This list is exhaustive, because \c escape sequences cannot be nested, and character classes cannot be nested directly. The choice of numbers is such as to optimize the most useful tests, and make transitions from one mode to another as simple as possible.

- Even modes mean that we are not directly in a character class. In this case, a left bracket appends 3 to the mode. In a character class, a right bracket changes the mode as $m \rightarrow (m - 15)/13$, truncated.
- Grouping, assertion, and anchors are allowed in non-positive even modes (0, -2, -6), and do not change the mode. Otherwise, they trigger an error.
- A left bracket is special in even modes, appending 3 to the mode; in those modes, quantifiers and the dot are recognized, and the right bracket is normal. In odd modes (within classes), the left bracket is normal, but the right bracket ends the class, changing the mode from m to $(m - 15)/13$, truncated; also, ranges are recognized.
- In non-negative modes, left and right braces are normal. In negative modes, however, left braces trigger a warning; right braces end the control sequence, going from -2 to 0 or -6 to 3, with error recovery for odd modes.
- Properties (such as the `\d` character class) can appear in any mode.

`_regex_if_in_class:TF` Test whether we are directly in a character class (at the innermost level of nesting). There, many escape sequences are not recognized, and special characters are normal. Also, for every raw character, we must look ahead for a possible raw dash.

```

24636 \cs_new:Npn \_regex_if_in_class:TF
24637   {
24638     \if_int_odd:w \l__regex_mode_int
24639       \exp_after:wN \use_i:nn
24640     \else:
24641       \exp_after:wN \use_ii:nn
24642     \fi:
24643   }

```

(End definition for `_regex_if_in_class:TF`.)

`_regex_if_in_cs:TF` Right braces are special only directly inside control sequences (at the inner-most level of nesting, not counting groups).

```

24644 \cs_new:Npn \_regex_if_in_cs:TF
24645   {
24646     \if_int_odd:w \l__regex_mode_int
24647       \exp_after:wN \use_ii:nn
24648     \else:
24649       \if_int_compare:w \l__regex_mode_int < \c__regex_outer_mode_int
24650         \exp_after:wN \exp_after:wN \exp_after:wN \use_i:nn
24651       \else:
24652         \exp_after:wN \exp_after:wN \exp_after:wN \use_ii:nn
24653       \fi:
24654     \fi:
24655   }

```

(End definition for `_regex_if_in_cs:TF`.)

`_regex_if_in_class_or_catcode:TF` Assertions are only allowed in modes 0, -2, and -6, *i.e.*, even, non-positive modes.

```

24656 \cs_new:Npn \_regex_if_in_class_or_catcode:TF
24657   {
24658     \if_int_odd:w \l__regex_mode_int

```



```

24659     \exp_after:wN \use_i:nn
24660 \else:
24661     \if_int_compare:w \l__regex_mode_int > \c__regex_outer_mode_int
24662     \exp_after:wN \exp_after:wN \exp_after:wN \use_i:nn
24663 \else:
24664     \exp_after:wN \exp_after:wN \exp_after:wN \use_ii:nn
24665 \fi:
24666 \fi:
24667 }

```

(End definition for __regex_if_in_class_or_catcode:TF.)

__regex_if_within_catcode:TF This test takes the true branch if we are in a catcode test, either immediately following it (modes 2 and 6) or in a class on which it applies (modes 23 and 63). This is used to tweak how left brackets behave in modes 2 and 6.

```

24668 \cs_new:Npn \__regex_if_within_catcode:TF
24669 {
24670     \if_int_compare:w \l__regex_mode_int > \c__regex_outer_mode_int
24671     \exp_after:wN \use_i:nn
24672 \else:
24673     \exp_after:wN \use_ii:nn
24674 \fi:
24675 }

```

(End definition for __regex_if_within_catcode:TF.)

__regex_chk_c_allowed:T The \c escape sequence is only allowed in modes 0 and 3, *i.e.*, not within any other \c escape sequence.

```

24676 \cs_new_protected:Npn \__regex_chk_c_allowed:T
24677 {
24678     \if_int_compare:w \l__regex_mode_int = \c__regex_outer_mode_int
24679     \exp_after:wN \use:n
24680 \else:
24681     \if_int_compare:w \l__regex_mode_int = \c__regex_class_mode_int
24682     \exp_after:wN \exp_after:wN \exp_after:wN \use:n
24683 \else:
24684     \__kernel_msg_error:nn { kernel } { c-bad-mode }
24685     \exp_after:wN \exp_after:wN \exp_after:wN \use_none:n
24686 \fi:
24687 \fi:
24688 }

```

(End definition for __regex_chk_c_allowed:T.)

__regex_mode_quit_c: This function changes the mode as it is needed just after a catcode test.

```

24689 \cs_new_protected:Npn \__regex_mode_quit_c:
24690 {
24691     \if_int_compare:w \l__regex_mode_int = \c__regex_catcode_mode_int
24692     \int_set_eq:NN \l__regex_mode_int \c__regex_outer_mode_int
24693 \else:
24694     \if_int_compare:w \l__regex_mode_int =
24695     \c__regex_catcode_in_class_mode_int
24696     \int_set_eq:NN \l__regex_mode_int \c__regex_class_mode_int
24697 \fi:
24698 \fi:
24699 }

```

(End definition for _regex_mode_quit_c:.)

41.3.4 Framework

_regex_compile:w Used when compiling a user regex or a regex for the \c{...} escape sequence within
_regex_compile_end: another regex. Start building a token list within a group (with x-expansion at the outset),
and set a few variables (group level, catcodes), then start the first branch. At the end,
make sure there are no dangling classes nor groups, close the last branch: we are done
building \l_regex_internal_regex.

```

24700 \cs_new_protected:Npn \_regex_compile:w
24701   {
24702     \group_begin:
24703     \tl_build_begin:N \l\_regex_build_tl
24704     \int_zero:N \l\_regex_group_level_int
24705     \int_set_eq:NN \l\_regex_default_catcodes_int
24706       \c\_regex_all_catcodes_int
24707     \int_set_eq:NN \l\_regex_catcodes_int \l\_regex_default_catcodes_int
24708     \cs_set:Npn \_regex_item_equal:n { \_regex_item_caseful_equal:n }
24709     \cs_set:Npn \_regex_item_range:nn { \_regex_item_caseful_range:nn }
24710     \tl_build_put_right:Nn \l\_regex_build_tl
24711       { \_regex_branch:n { \if_false: } \fi: }
24712   }
24713 \cs_new_protected:Npn \_regex_compile_end:
24714   {
24715     \_regex_if_in_class:TF
24716     {
24717       \_kernel_msg_error:nn { kernel } { missing-rbrack }
24718       \use:c { \_regex_compile_: }
24719       \prg_do_nothing: \prg_do_nothing:
24720     }
24721     { }
24722     \if_int_compare:w \l\_regex_group_level_int > 0 \exp_stop_f:
24723     \_kernel_msg_error:nnx { kernel } { missing-rparen }
24724     { \int_use:N \l\_regex_group_level_int }
24725     \prg_replicate:nn
24726       { \l\_regex_group_level_int }
24727     {
24728       \tl_build_put_right:Nn \l\_regex_build_tl
24729       {
24730         \if_false: { \fi: }
24731         \if_false: { \fi: } { 1 } { 0 } \c_true_bool
24732       }
24733       \tl_build_end:N \l\_regex_build_tl
24734       \exp_args:NNNo
24735       \group_end:
24736       \tl_build_put_right:Nn \l\_regex_build_tl
24737         { \l\_regex_build_tl }
24738     }
24739     \fi:
24740     \tl_build_put_right:Nn \l\_regex_build_tl { \if_false: { \fi: } }
24741     \tl_build_end:N \l\_regex_build_tl
24742     \exp_args:NNNx
24743     \group_end:

```

```

24744 \tl_set:Nn \l__regex_internal_regex { \l__regex_build_tl }
24745 }

```

(End definition for __regex_compile:w and __regex_compile_end:.)

__regex_compile:n The compilation is done between __regex_compile:w and __regex_compile_end:, starting in mode 0. Then __regex_escape_use:nnnn distinguishes special characters, escaped alphanumerics, and raw characters, interpreting \a, \x and other sequences. The 4 trailing \prg_do_nothing: are needed because some functions defined later look up to 4 tokens ahead. Before ending, make sure that any \c{...} is properly closed. No need to check that brackets are closed properly since __regex_compile_end: does that. However, catch the case of a trailing \cL construction.

```

24746 \cs_new_protected:Npn \__regex_compile:n #1
24747 {
24748   \__regex_compile:w
24749   \__regex_standard_escapechar:
24750   \int_set_eq:NN \l__regex_mode_int \c__regex_outer_mode_int
24751   \__regex_escape_use:nnnn
24752   {
24753     \__regex_char_if_special:NTF ##1
24754     \__regex_compile_special:N \__regex_compile_raw:N ##1
24755   }
24756   {
24757     \__regex_char_if_alphanumeric:NTF ##1
24758     \__regex_compile_escaped:N \__regex_compile_raw:N ##1
24759   }
24760   { \__regex_compile_raw:N ##1 }
24761   { #1 }
24762   \prg_do_nothing: \prg_do_nothing:
24763   \prg_do_nothing: \prg_do_nothing:
24764   \int_compare:nNnT \l__regex_mode_int = \c__regex_catcode_mode_int
24765   { \__kernel_msg_error:nn { kernel } { c-trailing } }
24766   \int_compare:nNnT \l__regex_mode_int < \c__regex_outer_mode_int
24767   {
24768     \__kernel_msg_error:nn { kernel } { c-missing-rbrace }
24769     \__regex_compile_end_cs:
24770     \prg_do_nothing: \prg_do_nothing:
24771     \prg_do_nothing: \prg_do_nothing:
24772   }
24773   \__regex_compile_end:
24774 }

```

(End definition for __regex_compile:n.)

__regex_compile_escaped:N If the special character or escaped alphanumeric has a particular meaning in regexes, the corresponding function is used. Otherwise, it is interpreted as a raw character. We distinguish special characters from escaped alphanumeric characters because they behave differently when appearing as an end-point of a range.

```

24775 \cs_new_protected:Npn \__regex_compile_special:N #1
24776 {
24777   \cs_if_exist_use:cF { __regex_compile_#1: }
24778   { \__regex_compile_raw:N #1 }
24779 }
24780 \cs_new_protected:Npn \__regex_compile_escaped:N #1

```

```

24781 {
24782   \cs_if_exist_use:cF { __regex_compile_/#1: }
24783   { \__regex_compile_raw:N #1 }
24784 }

```

(End definition for __regex_compile_escaped:N and __regex_compile_special:N.)

__regex_compile_one:n This is used after finding one “test”, such as \d, or a raw character. If that followed a catcode test (e.g., \cL), then restore the mode. If we are not in a class, then the test is “standalone”, and we need to add __regex_class:NnnnN and search for quantifiers. In any case, insert the test, possibly together with a catcode test if appropriate.

```

24785 \cs_new_protected:Npn \__regex_compile_one:n #1
24786 {
24787   \__regex_mode_quit_c:
24788   \__regex_if_in_class:TF { }
24789   {
24790     \tl_build_put_right:Nn \l__regex_build_tl
24791     { \__regex_class:NnnnN \c_true_bool { \if_false: } \fi: }
24792   }
24793   \tl_build_put_right:Nx \l__regex_build_tl
24794   {
24795     \if_int_compare:w \l__regex_catcodes_int <
24796     \c__regex_all_catcodes_int
24797     \__regex_item_catcode:nT { \int_use:N \l__regex_catcodes_int }
24798     { \exp_not:N \exp_not:n {#1} }
24799     \else:
24800     \exp_not:N \exp_not:n {#1}
24801     \fi:
24802   }
24803   \int_set_eq:NN \l__regex_catcodes_int \l__regex_default_catcodes_int
24804   \__regex_if_in_class:TF { } { \__regex_compile_quantifier:w }
24805 }

```

(End definition for __regex_compile_one:n.)

__regex_compile_abort_tokens:n This function places the collected tokens back in the input stream, each as a raw character.
 __regex_compile_abort_tokens:x Spaces are not preserved.

```

24806 \cs_new_protected:Npn \__regex_compile_abort_tokens:n #1
24807 {
24808   \use:x
24809   {
24810     \exp_args:No \tl_map_function:nN { \tl_to_str:n {#1} }
24811     \__regex_compile_raw:N
24812   }
24813 }
24814 \cs_generate_variant:Nn \__regex_compile_abort_tokens:n { x }

```

(End definition for __regex_compile_abort_tokens:n.)

41.3.5 Quantifiers

__regex_compile_quantifier:w This looks ahead and finds any quantifier (special character equal to either of ?+*{).

```

24815 \cs_new_protected:Npn \__regex_compile_quantifier:w #1#2
24816 {

```

```

24817 \token_if_eq_meaning:NNTF #1 \__regex_compile_special:N
24818 {
24819     \cs_if_exist_use:cF { __regex_compile_quantifier_#2:w }
24820     { \__regex_compile_quantifier_none: #1 #2 }
24821 }
24822 { \__regex_compile_quantifier_none: #1 #2 }
24823 }

```

(End definition for __regex_compile_quantifier:w.)

__regex_compile_quantifier_none: Those functions are called whenever there is no quantifier, or a braced construction is invalid (equivalent to no quantifier, and whatever characters were grabbed are left raw).
 __regex_compile_quantifier_abort:xNN

```

24824 \cs_new_protected:Npn \__regex_compile_quantifier_none:
24825 {
24826     \tl_build_put_right:Nn \l__regex_build_tl
24827     { \if_false: { \fi: } { 1 } { 0 } \c_false_bool }
24828 }
24829 \cs_new_protected:Npn \__regex_compile_quantifier_abort:xNN #1#2#3
24830 {
24831     \__regex_compile_quantifier_none:
24832     \__kernel_msg_warning:nxxx { kernel } { invalid-quantifier } {#1} {#3}
24833     \__regex_compile_abort_tokens:x {#1}
24834     #2 #3
24835 }

```

(End definition for __regex_compile_quantifier_none: and __regex_compile_quantifier_abort:xNN.)

__regex_compile_quantifier_lazyness:nnNN Once the “main” quantifier (?, *, + or a braced construction) is found, we check whether it is lazy (followed by a question mark). We then add to the compiled regex a closing brace (ending __regex_class:NnnnN and friends), the start-point of the range, its end-point, and a boolean, true for lazy and false for greedy operators.

```

24836 \cs_new_protected:Npn \__regex_compile_quantifier_lazyness:nnNN #1#2#3#4
24837 {
24838     \__regex_two_if_eq:NNNTF #3 #4 \__regex_compile_special:N ?
24839     {
24840         \tl_build_put_right:Nn \l__regex_build_tl
24841         { \if_false: { \fi: } { #1 } { #2 } \c_true_bool }
24842     }
24843     {
24844         \tl_build_put_right:Nn \l__regex_build_tl
24845         { \if_false: { \fi: } { #1 } { #2 } \c_false_bool }
24846         #3 #4
24847     }
24848 }

```

(End definition for __regex_compile_quantifier_lazyness:nnNN.)

__regex_compile_quantifier_?:w For each “basic” quantifier, ?, *, +, feed the correct arguments to __regex_compile_quantifier_lazyness:nnNN, -1 means that there is no upper bound on the number of repetitions.
 __regex_compile_quantifier_*:w
 __regex_compile_quantifier_+:w

```

24849 \cs_new_protected:cpn { __regex_compile_quantifier_?:w }
24850 { \__regex_compile_quantifier_lazyness:nnNN { 0 } { 1 } }
24851 \cs_new_protected:cpn { __regex_compile_quantifier_*:w }
24852 { \__regex_compile_quantifier_lazyness:nnNN { 0 } { -1 } }

```

```

24853 \cs_new_protected:cpn { __regex_compile_quantifier_+:w }
24854 { \__regex_compile_quantifier_lazyness:nnNN { 1 } { -1 } }

```

(End definition for `__regex_compile_quantifier_?:w`, `__regex_compile_quantifier_+:w`, and `__regex_compile_quantifier_+:w`.)

```

\__regex_compile_quantifier_{:w
\__regex_compile_quantifier_braced_auxi:w
\__regex_compile_quantifier_braced_auxii:w
\__regex_compile_quantifier_braced_auxiii:w

```

Three possible syntaxes: $\{\langle int \rangle\}$, $\{\langle int \rangle, \}$, or $\{\langle int \rangle, \langle int \rangle\}$. Any other syntax causes us to abort and put whatever we collected back in the input stream, as `raw` characters, including the opening brace. Grab a number into `\l__regex_internal_a_int`. If the number is followed by a right brace, the range is $[a, a]$. If followed by a comma, grab one more number, and call the `_ii` or `_iii` auxiliary. Those auxiliaries check for a closing brace, leading to the range $[a, \infty]$ or $[a, b]$, encoded as $\{a\}\{-1\}$ and $\{a\}\{b-a\}$.

```

24855 \cs_new_protected:cpn { __regex_compile_quantifier_ \c_left_brace_str :w }
24856 {
24857   \__regex_get_digits:NTFw \l__regex_internal_a_int
24858   { \__regex_compile_quantifier_braced_auxi:w }
24859   { \__regex_compile_quantifier_abort:xNN { \c_left_brace_str } }
24860 }
24861 \cs_new_protected:Npn \__regex_compile_quantifier_braced_auxi:w #1#2
24862 {
24863   \str_case_e:nnF { #1 #2 }
24864   {
24865     { \__regex_compile_special:N \c_right_brace_str }
24866     {
24867       \exp_args:No \__regex_compile_quantifier_lazyness:nnNN
24868       { \int_use:N \l__regex_internal_a_int } { 0 }
24869     }
24870     { \__regex_compile_special:N , }
24871     {
24872       \__regex_get_digits:NTFw \l__regex_internal_b_int
24873       { \__regex_compile_quantifier_braced_auxiii:w }
24874       { \__regex_compile_quantifier_braced_auxii:w }
24875     }
24876   }
24877   {
24878     \__regex_compile_quantifier_abort:xNN
24879     { \c_left_brace_str \int_use:N \l__regex_internal_a_int }
24880     #1 #2
24881   }
24882 }
24883 \cs_new_protected:Npn \__regex_compile_quantifier_braced_auxii:w #1#2
24884 {
24885   \__regex_two_if_eq:NNNTF #1 #2 \__regex_compile_special:N \c_right_brace_str
24886   {
24887     \exp_args:No \__regex_compile_quantifier_lazyness:nnNN
24888     { \int_use:N \l__regex_internal_a_int } { -1 }
24889   }
24890   {
24891     \__regex_compile_quantifier_abort:xNN
24892     { \c_left_brace_str \int_use:N \l__regex_internal_a_int , }
24893     #1 #2
24894   }
24895 }
24896 \cs_new_protected:Npn \__regex_compile_quantifier_braced_auxiii:w #1#2

```

```

24897 {
24898   \_regex_two_if_eq:NNNTF #1 #2 \_regex_compile_special:N \c_right_brace_str
24899   {
24900     \if_int_compare:w \l__regex_internal_a_int >
24901     \l__regex_internal_b_int
24902     \_kernel_msg_error:nnxx { kernel } { backwards-quantifier }
24903     { \int_use:N \l__regex_internal_a_int }
24904     { \int_use:N \l__regex_internal_b_int }
24905     \int_zero:N \l__regex_internal_b_int
24906   \else:
24907     \int_sub:Nn \l__regex_internal_b_int \l__regex_internal_a_int
24908   \fi:
24909   \exp_args:Noo \_regex_compile_quantifier_lazyness:nnNN
24910   { \int_use:N \l__regex_internal_a_int }
24911   { \int_use:N \l__regex_internal_b_int }
24912 }
24913 {
24914   \_regex_compile_quantifier_abort:xNN
24915   {
24916     \c_left_brace_str
24917     \int_use:N \l__regex_internal_a_int ,
24918     \int_use:N \l__regex_internal_b_int
24919   }
24920   #1 #2
24921 }
24922 }

```

(End definition for _regex_compile_quantifier_{:w and others.)

41.3.6 Raw characters

_regex_compile_raw_error:N Within character classes, and following catcode tests, some escaped alphanumeric sequences such as \b do not have any meaning. They are replaced by a raw character, after spitting out an error.

```

24923 \cs_new_protected:Npn \_regex_compile_raw_error:N #1
24924 {
24925   \_kernel_msg_error:nnx { kernel } { bad-escape } {#1}
24926   \_regex_compile_raw:N #1
24927 }

```

(End definition for _regex_compile_raw_error:N.)

_regex_compile_raw:N If we are in a character class and the next character is an unescaped dash, this denotes a range. Otherwise, the current character #1 matches itself.

```

24928 \cs_new_protected:Npn \_regex_compile_raw:N #1#2#3
24929 {
24930   \_regex_if_in_class:TF
24931   {
24932     \_regex_two_if_eq:NNNTF #2 #3 \_regex_compile_special:N -
24933     { \_regex_compile_range:Nw #1 }
24934     {
24935       \_regex_compile_one:n
24936       { \_regex_item_equal:n { \int_value:w '#1 } }
24937       #2 #3

```

```

24938     }
24939   }
24940   {
24941     \_regex_compile_one:n
24942     { \_regex_item_equal:n { \int_value:w '#1 } }
24943     #2 #3
24944   }
24945 }

```

(End definition for _regex_compile_raw:N.)

_regex_compile_range:Nw
_regex_if_end_range:NNTF

We have just read a raw character followed by a dash; this should be followed by an end-point for the range. Valid end-points are: any raw character; any special character, except a right bracket. In particular, escaped characters are forbidden.

```

24946 \prg_new_protected_conditional:Npnn \_regex_if_end_range:NN #1#2 { TF }
24947 {
24948   \if_meaning:w \_regex_compile_raw:N #1
24949   \prg_return_true:
24950   \else:
24951     \if_meaning:w \_regex_compile_special:N #1
24952     \if_charcode:w ] #2
24953     \prg_return_false:
24954     \else:
24955       \prg_return_true:
24956       \fi:
24957     \else:
24958       \prg_return_false:
24959       \fi:
24960   \fi:
24961 }
24962 \cs_new_protected:Npn \_regex_compile_range:Nw #1#2#3
24963 {
24964   \_regex_if_end_range:NNTF #2 #3
24965   {
24966     \if_int_compare:w '#1 > '#3 \exp_stop_f:
24967     \_kernel_msg_error:nxxx { kernel } { range-backwards } {#1} {#3}
24968     \else:
24969       \tl_build_put_right:Nx \l__regex_build_tl
24970       {
24971         \if_int_compare:w '#1 = '#3 \exp_stop_f:
24972         \_regex_item_equal:n
24973         \else:
24974           \_regex_item_range:nn { \int_value:w '#1 }
24975         \fi:
24976         { \int_value:w '#3 }
24977       }
24978     \fi:
24979   }
24980   {
24981     \_kernel_msg_warning:nxxx { kernel } { range-missing-end }
24982     {#1} { \c_backslash_str #3 }
24983     \tl_build_put_right:Nx \l__regex_build_tl
24984     {
24985       \_regex_item_equal:n { \int_value:w '#1 \exp_stop_f: }

```



```

24986         \_regex_item_equal:n { \int_value:w ‘- \exp_stop_f: }
24987     }
24988     #2#3
24989 }
24990 }

```

(End definition for _regex_compile_range:Nw and _regex_if_end_range:NNTF.)

41.3.7 Character properties

_regex_compile_.: In a class, the dot has no special meaning. Outside, insert _regex_prop_., which matches any character or control sequence, and refuses -2 (end-marker).

```

24991 \cs_new_protected:cpx { \_regex_compile_.: }
24992 {
24993     \exp_not:N \_regex_if_in_class:TF
24994     { \_regex_compile_raw:N . }
24995     { \_regex_compile_one:n \exp_not:c { \_regex_prop_.: } }
24996 }
24997 \cs_new_protected:cpn { \_regex_prop_.: }
24998 {
24999     \if_int_compare:w \l__regex_curr_char_int > - 2 \exp_stop_f:
25000     \exp_after:wN \_regex_break_true:w
25001     \fi:
25002 }

```

(End definition for _regex_compile_.: and _regex_prop_.:)

_regex_compile_/d: The constants _regex_prop_d:, etc. hold a list of tests which match the corresponding character class, and jump to the _regex_break_point:TF marker. As for a normal character, we check for quantifiers.

```

\_regex_compile_/d: 25003 \cs_set_protected:Npn \_regex_tmp:w #1#2
\_regex_compile_/D: 25004 {
\_regex_compile_/h: 25005     \cs_new_protected:cpx { \_regex_compile_/#1: }
\_regex_compile_/S: 25006     { \_regex_compile_one:n \exp_not:c { \_regex_prop_#1: } }
\_regex_compile_/v: 25007     \cs_new_protected:cpx { \_regex_compile_/#2: }
\_regex_compile_/V: 25008     {
\_regex_compile_/w: 25009         \_regex_compile_one:n
\_regex_compile_/W: 25010         { \_regex_item_reverse:n \exp_not:c { \_regex_prop_#1: } }
\_regex_compile_/N: 25011     }
25012 }
25013 \_regex_tmp:w d D
25014 \_regex_tmp:w h H
25015 \_regex_tmp:w s S
25016 \_regex_tmp:w v V
25017 \_regex_tmp:w w W
25018 \cs_new_protected:cpn { \_regex_compile_/N: }
25019 { \_regex_compile_one:n \_regex_prop_N: }

```

(End definition for _regex_compile_/d: and others.)

41.3.8 Anchoring and simple assertions

`_regex_compile_anchor_letter:NNN` In modes where assertions are forbidden, anchors such as `\A` produce an error (`\A` is invalid in classes); otherwise they add an `_regex_assertion:Nn` test as appropriate (the only negative assertion is `\B`). The test functions are defined later. The implementation for `$` and `^` is only different from `\A` etc because these are valid in a class.

```

25020 \cs_new_protected:Npn \_regex_compile_anchor_letter:NNN #1#2#3
25021 {
25022   \_regex_if_in_class_or_catcode:TF { \_regex_compile_raw_error:N #1 }
25023   {
25024     \tl_build_put_right:Nn \l__regex_build_tl
25025     { \_regex_assertion:Nn #2 {#3} }
25026   }
25027 }
25028 \cs_new_protected:cpn { __regex_compile_/A: }
25029 { \_regex_compile_anchor_letter:NNN A \c_true_bool \_regex_A_test: }
25030 \cs_new_protected:cpn { __regex_compile_/G: }
25031 { \_regex_compile_anchor_letter:NNN G \c_true_bool \_regex_G_test: }
25032 \cs_new_protected:cpn { __regex_compile_/Z: }
25033 { \_regex_compile_anchor_letter:NNN Z \c_true_bool \_regex_Z_test: }
25034 \cs_new_protected:cpn { __regex_compile_/z: }
25035 { \_regex_compile_anchor_letter:NNN z \c_true_bool \_regex_Z_test: }
25036 \cs_new_protected:cpn { __regex_compile_/b: }
25037 { \_regex_compile_anchor_letter:NNN b \c_true_bool \_regex_b_test: }
25038 \cs_new_protected:cpn { __regex_compile_/B: }
25039 { \_regex_compile_anchor_letter:NNN B \c_false_bool \_regex_b_test: }
25040 \cs_set_protected:Npn \_regex_tmp:w #1#2
25041 {
25042   \cs_new_protected:cpn { __regex_compile_#1: }
25043   {
25044     \_regex_if_in_class_or_catcode:TF { \_regex_compile_raw:N #1 }
25045     {
25046       \tl_build_put_right:Nn \l__regex_build_tl
25047       { \_regex_assertion:Nn \c_true_bool {#2} }
25048     }
25049   }
25050 }
25051 \exp_args:Nx \_regex_tmp:w { \iow_char:N \^ } { \_regex_A_test: }
25052 \exp_args:Nx \_regex_tmp:w { \iow_char:N \$ } { \_regex_Z_test: }

```

(End definition for `_regex_compile_anchor_letter:NNN` and others.)

41.3.9 Character classes

`_regex_compile_]:` Outside a class, right brackets have no meaning. In a class, change the mode ($m \rightarrow (m - 15)/13$, truncated) to reflect the fact that we are leaving the class. Look for quantifiers, unless we are still in a class after leaving one (the case of `[... \cL[...]`...). quantifiers.

```

25053 \cs_new_protected:cpn { __regex_compile_]: }
25054 {
25055   \_regex_if_in_class:TF
25056   {
25057     \if_int_compare:w \l__regex_mode_int >
25058     \c__regex_catcode_in_class_mode_int
25059     \tl_build_put_right:Nn \l__regex_build_tl { \if_false: { \fi: } }

```

```

25060     \fi:
25061     \tex_advance:D \l__regex_mode_int - 15 \exp_stop_f:
25062     \tex_divide:D \l__regex_mode_int 13 \exp_stop_f:
25063     \if_int_odd:w \l__regex_mode_int \else:
25064         \exp_after:wN \__regex_compile_quantifier:w
25065     \fi:
25066 }
25067 { \__regex_compile_raw:N ] }
25068 }

```

(End definition for __regex_compile_[:.))

__regex_compile_[: In a class, left brackets might introduce a POSIX character class, or mean nothing. Immediately following \c<category>, we must insert the appropriate catcode test, then parse the class; we pre-expand the catcode as an optimization. Otherwise (modes 0, -2 and -6) just parse the class. The mode is updated later.

```

25069 \cs_new_protected:cpn { \__regex_compile_[: }
25070 {
25071     \__regex_if_in_class:TF
25072     { \__regex_compile_class_posix_test:w }
25073     {
25074         \__regex_if_within_catcode:TF
25075         {
25076             \exp_after:wN \__regex_compile_class_catcode:w
25077             \int_use:N \l__regex_catcodes_int ;
25078         }
25079         { \__regex_compile_class_normal:w }
25080     }
25081 }

```

(End definition for __regex_compile_[:.))

__regex_compile_class_normal:w In the “normal” case, we insert __regex_class:NnnnN <boolean> in the compiled code. The <boolean> is true for positive classes, and false for negative classes, characterized by a leading ~. The auxiliary __regex_compile_class:TFNN also checks for a leading] which has a special meaning.

```

25082 \cs_new_protected:Npn \__regex_compile_class_normal:w
25083 {
25084     \__regex_compile_class:TFNN
25085     { \__regex_class:NnnnN \c_true_bool }
25086     { \__regex_class:NnnnN \c_false_bool }
25087 }

```

(End definition for __regex_compile_class_normal:w.))

__regex_compile_class_catcode:w This function is called for a left bracket in modes 2 or 6 (catcode test, and catcode test within a class). In mode 2 the whole construction needs to be put in a class (like single character). Then determine if the class is positive or negative, inserting __regex_item_catcode:nT or the reverse variant as appropriate, each with the current catcodes bitmap #1 as an argument, and reset the catcodes.

```

25088 \cs_new_protected:Npn \__regex_compile_class_catcode:w #1;
25089 {
25090     \if_int_compare:w \l__regex_mode_int = \c__regex_catcode_mode_int
25091         \tl_build_put_right:Nn \l__regex_build_tl

```

```

25092         { \__regex_class:NnnnN \c_true_bool { \if_false: } \fi: }
25093     \fi:
25094     \int_set_eq:NN \l__regex_catcodes_int \l__regex_default_catcodes_int
25095     \__regex_compile_class:TFNN
25096     { \__regex_item_catcode:nT {#1} }
25097     { \__regex_item_catcode_reverse:nT {#1} }
25098 }

```

(End definition for __regex_compile_class_catcode:w.)

__regex_compile_class:TFNN If the first character is ^, then the class is negative (use #2), otherwise it is positive (use #1). If the next character is a right bracket, then it should be changed to a raw one.

```

25099 \cs_new_protected:Npn \__regex_compile_class:TFNN #1#2#3#4
25100 {
25101     \l__regex_mode_int = \int_value:w \l__regex_mode_int 3 \exp_stop_f:
25102     \__regex_two_if_eq:NNNTF #3 #4 \__regex_compile_special:N ^
25103     {
25104         \tl_build_put_right:Nn \l__regex_build_tl { #2 { \if_false: } \fi: }
25105         \__regex_compile_class:NN
25106     }
25107     {
25108         \tl_build_put_right:Nn \l__regex_build_tl { #1 { \if_false: } \fi: }
25109         \__regex_compile_class:NN #3 #4
25110     }
25111 }
25112 \cs_new_protected:Npn \__regex_compile_class:NN #1#2
25113 {
25114     \token_if_eq_charcode:NNTF #2 ]
25115     { \__regex_compile_raw:N #2 }
25116     { #1 #2 }
25117 }

```

(End definition for __regex_compile_class:TFNN and __regex_compile_class:NN.)

__regex_compile_class_posix_test:w Here we check for a syntax such as [:alpha:]. We also detect [= and [. which have a meaning in POSIX regular expressions, but are not implemented in l3regex. In case we see [:, grab raw characters until hopefully reaching :]. If that's missing, or the POSIX class is unknown, abort. If all is right, add the test to the current class, with an extra __regex_item_reverse:n for negative classes.

```

25118 \cs_new_protected:Npn \__regex_compile_class_posix_test:w #1#2
25119 {
25120     \token_if_eq_meaning:NNT \__regex_compile_special:N #1
25121     {
25122         \str_case:nn { #2 }
25123         {
25124             : { \__regex_compile_class_posix:NNNNw }
25125             = {
25126                 \__kernel_msg_warning:nnx { kernel }
25127                 { posix-unsupported } { = }
25128             }
25129             . {
25130                 \__kernel_msg_warning:nnx { kernel }
25131                 { posix-unsupported } { . }
25132             }

```

```

25133     }
25134 }
25135 \__regex_compile_raw:N [ #1 #2
25136 }
25137 \cs_new_protected:Npn \__regex_compile_class_posix:NNNNw #1#2#3#4#5#6
25138 {
25139     \__regex_two_if_eq:NNNTF #5 #6 \__regex_compile_special:N ^
25140     {
25141         \bool_set_false:N \l__regex_internal_bool
25142         \__kernel_tl_set:Nx \l__regex_internal_a_tl { \if_false: } \fi:
25143         \__regex_compile_class_posix_loop:w
25144     }
25145     {
25146         \bool_set_true:N \l__regex_internal_bool
25147         \__kernel_tl_set:Nx \l__regex_internal_a_tl { \if_false: } \fi:
25148         \__regex_compile_class_posix_loop:w #5 #6
25149     }
25150 }
25151 \cs_new:Npn \__regex_compile_class_posix_loop:w #1#2
25152 {
25153     \token_if_eq_meaning:NNTF \__regex_compile_raw:N #1
25154     { #2 \__regex_compile_class_posix_loop:w }
25155     { \if_false: { \fi: } \__regex_compile_class_posix_end:w #1 #2 }
25156 }
25157 \cs_new_protected:Npn \__regex_compile_class_posix_end:w #1#2#3#4
25158 {
25159     \__regex_two_if_eq:NNNTF #1 #2 \__regex_compile_special:N :
25160     { \__regex_two_if_eq:NNNTF #3 #4 \__regex_compile_special:N ] }
25161     { \use_ii:nn }
25162     {
25163         \cs_if_exist:cTF { __regex_posix_ \l__regex_internal_a_tl : }
25164         {
25165             \__regex_compile_one:n
25166             {
25167                 \bool_if:NF \l__regex_internal_bool \__regex_item_reverse:n
25168                 \exp_not:c { __regex_posix_ \l__regex_internal_a_tl : }
25169             }
25170         }
25171         {
25172             \__kernel_msg_warning:nxx { kernel } { posix-unknown }
25173             { \l__regex_internal_a_tl }
25174             \__regex_compile_abort_tokens:x
25175             {
25176                 [: \bool_if:NF \l__regex_internal_bool { ^ }
25177                 \l__regex_internal_a_tl :]
25178             }
25179         }
25180     }
25181     {
25182         \__kernel_msg_error:nxxx { kernel } { posix-missing-close }
25183         { [: \l__regex_internal_a_tl ] { #2 #4 }
25184         \__regex_compile_abort_tokens:x { [: \l__regex_internal_a_tl }
25185         #1 #2 #3 #4
25186     }

```

```
25187 }
```

(End definition for `_regex_compile_class_posix_test:w` and others.)

41.3.10 Groups and alternations

```
\_regex_compile_group_begin:N
\_regex_compile_group_end:
```

The contents of a regex group are turned into compiled code in `\l__regex_build_tl`, which ends up with items of the form `_regex_branch:n {⟨concatenation⟩}`. This construction is done using `\tl_build_...` functions within a `TeX` group, which automatically makes sure that options (case-sensitivity and default catcode) are reset at the end of the group. The argument `#1` is `_regex_group:nnnN` or a variant thereof. A small subtlety to support `\cL(abc)` as a shorthand for `(\cLa\cLb\cLc)`: exit any pending catcode test, save the category code at the start of the group as the default catcode for that group, and make sure that the catcode is restored to the default outside the group.

```
25188 \cs_new_protected:Npn \_regex_compile_group_begin:N #1
25189 {
25190   \tl_build_put_right:Nn \l__regex_build_tl { #1 { \if_false: } \fi: }
25191   \_regex_mode_quit_c:
25192   \group_begin:
25193     \tl_build_begin:N \l__regex_build_tl
25194     \int_set_eq:NN \l__regex_default_catcodes_int \l__regex_catcodes_int
25195     \int_incr:N \l__regex_group_level_int
25196     \tl_build_put_right:Nn \l__regex_build_tl
25197       { \_regex_branch:n { \if_false: } \fi: }
25198   }
25199 \cs_new_protected:Npn \_regex_compile_group_end:
25200 {
25201   \if_int_compare:w \l__regex_group_level_int > 0 \exp_stop_f:
25202     \tl_build_put_right:Nn \l__regex_build_tl { \if_false: { \fi: } }
25203     \tl_build_end:N \l__regex_build_tl
25204     \exp_args:NNNx
25205     \group_end:
25206     \tl_build_put_right:Nn \l__regex_build_tl { \l__regex_build_tl }
25207     \int_set_eq:NN \l__regex_catcodes_int \l__regex_default_catcodes_int
25208     \exp_after:wN \_regex_compile_quantifier:w
25209   \else:
25210     \__kernel_msg_warning:nn { kernel } { extra-rparen }
25211     \exp_after:wN \_regex_compile_raw:N \exp_after:wN )
25212   \fi:
25213 }
```

(End definition for `_regex_compile_group_begin:N` and `_regex_compile_group_end:.`)

`_regex_compile_(:` In a class, parentheses are not special. In a catcode test inside a class, a left parenthesis gives an error, to catch `[a\cL(bcd)e]`. Otherwise check for a `?`, denoting special groups, and run the code for the corresponding special group.

```
25214 \cs_new_protected:cpn { \_regex_compile_(: }
25215 {
25216   \_regex_if_in_class:TF { \_regex_compile_raw:N ( }
25217   {
25218     \if_int_compare:w \l__regex_mode_int =
25219       \c__regex_catcode_in_class_mode_int
```

```

25220         \_kernel_msg_error:nn { kernel } { c-lparen-in-class }
25221         \exp_after:wN \_regex_compile_raw:N \exp_after:wN (
25222         \else:
25223         \exp_after:wN \_regex_compile_lparen:w
25224         \fi:
25225     }
25226 }
25227 \cs_new_protected:Npn \_regex_compile_lparen:w #1#2#3#4
25228 {
25229     \_regex_two_if_eq:NNNTF #1 #2 \_regex_compile_special:N ?
25230     {
25231         \cs_if_exist_use:cF
25232         { \_regex_compile_special_group\_token_to_str:N #4 :w }
25233         {
25234             \_kernel_msg_warning:nxx { kernel } { special-group-unknown }
25235             { (? #4 }
25236             \_regex_compile_group_begin:N \_regex_group:nnnN
25237             \_regex_compile_raw:N ? #3 #4
25238         }
25239     }
25240     {
25241         \_regex_compile_group_begin:N \_regex_group:nnnN
25242         #1 #2 #3 #4
25243     }
25244 }

```

(End definition for _regex_compile(:))

_regex_compile_|: In a class, the pipe is not special. Otherwise, end the current branch and open another one.

```

25245 \cs_new_protected:cpn { \_regex_compile_|: }
25246 {
25247     \_regex_if_in_class:TF { \_regex_compile_raw:N | }
25248     {
25249         \tl_build_put_right:Nn \l\_regex_build_tl
25250         { \if_false: { \fi: } \_regex_branch:n { \if_false: } \fi: }
25251     }
25252 }

```

(End definition for _regex_compile_|:))

_regex_compile_): Within a class, parentheses are not special. Outside, close a group.

```

25253 \cs_new_protected:cpn { \_regex_compile_): }
25254 {
25255     \_regex_if_in_class:TF { \_regex_compile_raw:N ) }
25256     { \_regex_compile_group_end: }
25257 }

```

(End definition for _regex_compile_):))

_regex_compile_special_group::w Non-capturing, and resetting groups are easy to take care of during compilation; for those groups, the harder parts come when building.

```

25258 \cs_new_protected:cpn { \_regex_compile_special_group::w }
25259 { \_regex_compile_group_begin:N \_regex_group_no_capture:nnnN }
25260 \cs_new_protected:cpn { \_regex_compile_special_group_|:w }
25261 { \_regex_compile_group_begin:N \_regex_group_resetting:nnnN }

```

(End definition for `__regex_compile_special_group_i:w` and `__regex_compile_special_group_l:w`.)

`__regex_compile_special_group_i:w` The match can be made case-insensitive by setting the option with `(?i)`; the original behaviour is restored by `(?-i)`. This is the only supported option.

```

25262 \cs_new_protected:Npn __regex_compile_special_group_i:w #1#2
25263 {
25264   __regex_two_if_eq:NNNTF #1 #2 __regex_compile_special:N )
25265   {
25266     \cs_set:Npn __regex_item_equal:n
25267       { __regex_item_caseless_equal:n }
25268     \cs_set:Npn __regex_item_range:nn
25269       { __regex_item_caseless_range:nn }
25270   }
25271   {
25272     \__kernel_msg_warning:nxx { kernel } { unknown-option } { (?i #2 }
25273     __regex_compile_raw:N (
25274     __regex_compile_raw:N ?
25275     __regex_compile_raw:N i
25276     #1 #2
25277   }
25278 }
25279 \cs_new_protected:cpn { __regex_compile_special_group_~:w } #1#2#3#4
25280 {
25281   __regex_two_if_eq:NNNTF #1 #2 __regex_compile_raw:N i
25282   { __regex_two_if_eq:NNNTF #3 #4 __regex_compile_special:N ) }
25283   { \use_i:nn }
25284   {
25285     \cs_set:Npn __regex_item_equal:n
25286       { __regex_item_caseful_equal:n }
25287     \cs_set:Npn __regex_item_range:nn
25288       { __regex_item_caseful_range:nn }
25289   }
25290   {
25291     \__kernel_msg_warning:nxx { kernel } { unknown-option } { (?~#2#4 }
25292     __regex_compile_raw:N (
25293     __regex_compile_raw:N ?
25294     __regex_compile_raw:N -
25295     #1 #2 #3 #4
25296   }
25297 }

```

(End definition for `__regex_compile_special_group_i:w` and `__regex_compile_special_group_~:w`.)

41.3.11 Catcodes and csnames

`__regex_compile_/c:` The `\c` escape sequence can be followed by a capital letter representing a character category, by a left bracket which starts a list of categories, or by a brace group holding a regular expression for a control sequence name. Otherwise, raise an error.

```

25298 \cs_new_protected:cpn { __regex_compile_/c: }
25299 { __regex_chk_c_allowed:T { __regex_compile_c_test:NN } }
25300 \cs_new_protected:Npn __regex_compile_c_test:NN #1#2
25301 {
25302   \token_if_eq_meaning:NNTF #1 __regex_compile_raw:N
25303   {

```



```

25304     \int_if_exist:cTF { c__regex_catcode_#2_int }
25305     {
25306         \int_set_eq:Nc \l__regex_catcodes_int
25307         { c__regex_catcode_#2_int }
25308         \l__regex_mode_int
25309         = \if_case:w \l__regex_mode_int
25310           \c__regex_catcode_mode_int
25311           \else:
25312             \c__regex_catcode_in_class_mode_int
25313           \fi:
25314         \token_if_eq_charcode:NNT C #2 { \__regex_compile_c_C:NN }
25315     }
25316 }
25317 { \cs_if_exist_use:cF { __regex_compile_c_#2:w } }
25318 {
25319     \__kernel_msg_error:nmx { kernel } { c-missing-category } {#2}
25320     #1 #2
25321 }
25322 }

```

(End definition for __regex_compile_/c: and __regex_compile_c_test:NN.)

__regex_compile_c_C:NN If \cC is not followed by . or (...) then complain because that construction cannot match anything, except in cases like \cC[\c{...}], where it has no effect.

```

25323 \cs_new_protected:Npn \__regex_compile_c_C:NN #1#2
25324 {
25325     \token_if_eq_meaning:NNTF #1 \__regex_compile_special:N
25326     {
25327         \token_if_eq_charcode:NNTF #2 .
25328         { \use_none:n }
25329         { \token_if_eq_charcode:NNTF #2 ( ) % )
25330         }
25331         { \use:n }
25332         { \__kernel_msg_error:nnn { kernel } { c-C-invalid } {#2} }
25333         #1 #2
25334     }

```

(End definition for __regex_compile_c_C:NN.)

__regex_compile_c[:w] When encountering \c[, the task is to collect uppercase letters representing character categories. First check for ^ which negates the list of category codes.

```

\__regex_compile_c_lbrack_loop:NN
\__regex_compile_c_lbrack_add:N
\__regex_compile_c_lbrack_end:
25335 \cs_new_protected:cpn { __regex_compile_c[:w] } #1#2
25336 {
25337     \l__regex_mode_int
25338     = \if_case:w \l__regex_mode_int
25339       \c__regex_catcode_mode_int
25340       \else:
25341         \c__regex_catcode_in_class_mode_int
25342       \fi:
25343     \int_zero:N \l__regex_catcodes_int
25344     \__regex_two_if_eq:NNNTF #1 #2 \__regex_compile_special:N ^
25345     {
25346         \bool_set_false:N \l__regex_catcodes_bool
25347         \__regex_compile_c_lbrack_loop:NN

```

```

25348     }
25349     {
25350         \bool_set_true:N \l__regex_catcodes_bool
25351         \__regex_compile_c_lbrack_loop:NN
25352         #1 #2
25353     }
25354 }
25355 \cs_new_protected:Npn \__regex_compile_c_lbrack_loop:NN #1#2
25356 {
25357     \token_if_eq_meaning:NNTF #1 \__regex_compile_raw:N
25358     {
25359         \int_if_exist:cTF { c__regex_catcode_#2_int }
25360         {
25361             \exp_args:Nc \__regex_compile_c_lbrack_add:N
25362             { c__regex_catcode_#2_int }
25363             \__regex_compile_c_lbrack_loop:NN
25364         }
25365     }
25366     {
25367         \token_if_eq_charcode:NNTF #2 ]
25368         { \__regex_compile_c_lbrack_end: }
25369     }
25370     {
25371         \__kernel_msg_error:nxx { kernel } { c-missing-rbrack } {#2}
25372         \__regex_compile_c_lbrack_end:
25373         #1 #2
25374     }
25375 }
25376 \cs_new_protected:Npn \__regex_compile_c_lbrack_add:N #1
25377 {
25378     \if_int_odd:w \int_eval:n { \l__regex_catcodes_int / #1 } \exp_stop_f:
25379     \else:
25380         \int_add:Nn \l__regex_catcodes_int {#1}
25381     \fi:
25382 }
25383 \cs_new_protected:Npn \__regex_compile_c_lbrack_end:
25384 {
25385     \if_meaning:w \c_false_bool \l__regex_catcodes_bool
25386     \int_set:Nn \l__regex_catcodes_int
25387     { \c__regex_all_catcodes_int - \l__regex_catcodes_int }
25388     \fi:
25389 }

```

(End definition for __regex_compile_c[:w and others.)

__regex_compile_c_{: The case of a left brace is easy, based on what we have done so far: in a group, compile the regular expression, after changing the mode to forbid nesting \c. Additionally, disable submatch tracking since groups don't escape the scope of \c{...}.

```

25390 \cs_new_protected:cpn { __regex_compile_c_ \c_left_brace_str :w }
25391 {
25392     \__regex_compile:w
25393     \__regex_disable_submatches:
25394     \l__regex_mode_int
25395     = \if_case:w \l__regex_mode_int

```

```

25396         \c__regex_cs_mode_int
25397     \else:
25398         \c__regex_cs_in_class_mode_int
25399     \fi:
25400 }

```

(End definition for __regex_compile_c_{:})

__regex_compile_}: Non-escaped right braces are only special if they appear when compiling the regular expression for a csname, but not within a class: \c{[{}]} matches the control sequences \{ and \}. So, end compiling the inner regex (this closes any dangling class or group). Then insert the corresponding test in the outer regex. As an optimization, if the control sequence test simply consists of several explicit possibilities (branches) then use __regex_item_exact_cs:n with an argument consisting of all possibilities separated by \scan_stop:.

```

25401 \flag_new:n { __regex_cs }
25402 \cs_new_protected:cpn { __regex_compile_ \c_right_brace_str : }
25403 {
25404     \__regex_if_in_cs:TF
25405     { \__regex_compile_end_cs: }
25406     { \exp_after:wN \__regex_compile_raw:N \c_right_brace_str }
25407 }
25408 \cs_new_protected:Npn \__regex_compile_end_cs:
25409 {
25410     \__regex_compile_end:
25411     \flag_clear:n { __regex_cs }
25412     \__kernel_tl_set:Nx \l__regex_internal_a_tl
25413     {
25414         \exp_after:wN \__regex_compile_cs_aux:Nn \l__regex_internal_regex
25415         \q__regex_nil \q__regex_nil \q__regex_recursion_stop
25416     }
25417     \exp_args:Nx \__regex_compile_one:n
25418     {
25419         \flag_if_raised:nTF { __regex_cs }
25420         { \__regex_item_cs:n { \exp_not:o \l__regex_internal_regex } }
25421         {
25422             \__regex_item_exact_cs:n
25423             { \tl_tail:N \l__regex_internal_a_tl }
25424         }
25425     }
25426 }
25427 \cs_new:Npn \__regex_compile_cs_aux:Nn #1#2
25428 {
25429     \cs_if_eq:NNTF #1 \__regex_branch:n
25430     {
25431         \scan_stop:
25432         \__regex_compile_cs_aux:NNnnN #2
25433         \q__regex_nil \q__regex_nil \q__regex_nil
25434         \q__regex_nil \q__regex_nil \q__regex_nil \q__regex_recursion_stop
25435         \__regex_compile_cs_aux:Nn
25436     }
25437     {
25438         \__regex_quark_if_nil:NF #1 { \flag_raise_if_clear:n { __regex_cs } }
25439         \__regex_use_none_delimit_by_q_recursion_stop:w

```

```

25440     }
25441   }
25442   \cs_new:Npn \__regex_compile_cs_aux:NNnnnN #1#2#3#4#5#6
25443   {
25444     \bool_lazy_all:nTF
25445     {
25446       { \cs_if_eq_p:NN #1 \__regex_class:NnnnN }
25447       {#2}
25448       { \tl_if_head_eq_meaning_p:nN {#3} \__regex_item_caseful_equal:n }
25449       { \int_compare_p:nNn { \tl_count:n {#3} } = { 2 } }
25450       { \int_compare_p:nNn {#5} = { 0 } }
25451     }
25452     {
25453       \prg_replicate:nn {#4}
25454       { \char_generate:nn { \use_ii:nn #3 } {12} }
25455       \__regex_compile_cs_aux:NNnnnN
25456     }
25457     {
25458       \__regex_quark_if_nil:NF #1
25459       {
25460         \flag_raise_if_clear:n { __regex_cs }
25461         \__regex_use_i_delimit_by_q_recursion_stop:nw
25462       }
25463       \__regex_use_none_delimit_by_q_recursion_stop:w
25464     }
25465   }

```

(End definition for `__regex_compile_`: and others.)

41.3.12 Raw token lists with `\u`

`__regex_compile_/u:` The `\u` escape is invalid in classes and directly following a catcode test. Otherwise, it must be followed by a left brace. We then collect the characters for the argument of `\u` within an `x`-expanding assignment. In principle we could just wait to encounter a right brace, but this is unsafe: if the right brace was missing, then we would reach the end-markers of the regex, and continue, leading to obscure fatal errors. Instead, we only allow raw and special characters, and stop when encountering a special right brace, any escaped character, or the end-marker.

```

25466   \cs_new_protected:cpn { __regex_compile_/u: } #1#2
25467   {
25468     \__regex_if_in_class_or_catcode:TF
25469     { \__regex_compile_raw_error:N u #1 #2 }
25470     {
25471       \__regex_two_if_eq:NNNTF #1 #2 \__regex_compile_special:N \c_left_brace_str
25472       {
25473         \__kernel_tl_set:Nx \l__regex_internal_a_tl { \if_false: } \fi:
25474         \__regex_compile_u_loop:NN
25475       }
25476       {
25477         \__kernel_msg_error:nn { kernel } { u-missing-lbrace }
25478         \__regex_compile_raw:N u #1 #2
25479       }
25480     }
25481   }

```

```

25482 \cs_new:Npn \__regex_compile_u_loop:NN #1#2
25483 {
25484   \token_if_eq_meaning:NNTF #1 \__regex_compile_raw:N
25485   { #2 \__regex_compile_u_loop:NN }
25486   {
25487     \token_if_eq_meaning:NNTF #1 \__regex_compile_special:N
25488     {
25489       \exp_after:wN \token_if_eq_charcode:NNTF \c_right_brace_str #2
25490       { \if_false: { \fi: } \__regex_compile_u_end: }
25491       { #2 \__regex_compile_u_loop:NN }
25492     }
25493     {
25494       \if_false: { \fi: }
25495       \__kernel_msg_error:nxx { kernel } { u-missing-rbrace } {#2}
25496       \__regex_compile_u_end:
25497       #1 #2
25498     }
25499   }
25500 }

```

(End definition for __regex_compile_/u: and __regex_compile_u_loop:NN.)

__regex_compile_u_end: Once we have extracted the variable's name, we store the contents of that variable in \l__regex_internal_a_tl. The behaviour of \u then depends on whether we are within a \c{...} escape (in this case, the variable is turned to a string), or not.

```

25501 \cs_new_protected:Npn \__regex_compile_u_end:
25502 {
25503   \tl_set:Nv \l__regex_internal_a_tl { \l__regex_internal_a_tl }
25504   \if_int_compare:w \l__regex_mode_int = \c__regex_outer_mode_int
25505   \__regex_compile_u_not_cs:
25506   \else:
25507   \__regex_compile_u_in_cs:
25508   \fi:
25509 }

```

(End definition for __regex_compile_u_end:.)

__regex_compile_u_in_cs: When \u appears within a control sequence, we convert the variable to a string with escaped spaces. Then for each character insert a class matching exactly that character, once.

```

25510 \cs_new_protected:Npn \__regex_compile_u_in_cs:
25511 {
25512   \__kernel_tl_gset:Nx \g__regex_internal_tl
25513   {
25514     \exp_args:No \__kernel_str_to_other_fast:n
25515     { \l__regex_internal_a_tl }
25516   }
25517   \tl_build_put_right:Nx \l__regex_build_tl
25518   {
25519     \tl_map_function:NN \g__regex_internal_tl
25520     \__regex_compile_u_in_cs_aux:n
25521   }
25522 }
25523 \cs_new:Npn \__regex_compile_u_in_cs_aux:n #1

```

```

25524 {
25525   \_regex\_class:NnnnN \c\_true\_bool
25526   { \_regex\_item\_caseful\_equal:n { \int\_value:w '#1 } }
25527   { 1 } { 0 } \c\_false\_bool
25528 }

```

(End definition for _regex_compile_u_in_cs:.)

_regex_compile_u_not_cs: In mode 0, the \u escape adds one state to the NFA for each token in \l_regex_internal_a_tl. If a given *<token>* is a control sequence, then insert a string comparison test, otherwise, _regex_item_exact:nn which compares catcode and character code.

```

25529 \cs\_new\_protected:Npn \_regex\_compile\_u\_not\_cs:
25530 {
25531   \tl\_analysis\_map\_inline:Nn \l\_regex\_internal\_a\_tl
25532   {
25533     \tl\_build\_put\_right:Nx \l\_regex\_build\_tl
25534     {
25535       \_regex\_class:NnnnN \c\_true\_bool
25536       {
25537         \if\_int\_compare:w "##3 = 0 \exp\_stop\_f:
25538         \_regex\_item\_exact\_cs:n
25539         { \exp\_after:wN \cs\_to\_str:N ##1 }
25540         \else:
25541         \_regex\_item\_exact:nn { \int\_value:w "##3 } { ##2 }
25542         \fi:
25543       }
25544       { 1 } { 0 } \c\_false\_bool
25545     }
25546   }
25547 }

```

(End definition for _regex_compile_u_not_cs:.)

41.3.13 Other

_regex_compile_/K: The \K control sequence is currently the only “command”, which performs some action, rather than matching something. It is allowed in the same contexts as \b. At the compilation stage, we leave it as a single control sequence, defined later.

```

25548 \cs\_new\_protected:cpn { \_regex\_compile\_/K: }
25549 {
25550   \int\_compare:nNnTF \l\_regex\_mode\_int = \c\_regex\_outer\_mode\_int
25551   { \tl\_build\_put\_right:Nn \l\_regex\_build\_tl { \_regex\_command\_K: } }
25552   { \_regex\_compile\_raw\_error:N K }
25553 }

```

(End definition for _regex_compile_/K:.)

41.3.14 Showing regexes

_regex_show:N Within a group and within \tl_build_begin:N ... \tl_build_end:N we redefine all the function that can appear in a compiled regex, then run the regex. The result stored in \l_regex_internal_a_tl is then meant to be shown.

```

25554 \cs\_new\_protected:Npn \_regex\_show:N #1
25555 {

```

```

25556 \group_begin:
25557 \tl_build_begin:N \l__regex_build_tl
25558 \cs_set_protected:Npn \__regex_branch:n
25559 {
25560   \seq_pop_right:NN \l__regex_show_prefix_seq
25561   \l__regex_internal_a_tl
25562   \__regex_show_one:n { +-branch }
25563   \seq_put_right:No \l__regex_show_prefix_seq
25564   \l__regex_internal_a_tl
25565   \use:n
25566 }
25567 \cs_set_protected:Npn \__regex_group:nnnN
25568 { \__regex_show_group_aux:nnnnN { } }
25569 \cs_set_protected:Npn \__regex_group_no_capture:nnnN
25570 { \__regex_show_group_aux:nnnnN { ~(no~capture) } }
25571 \cs_set_protected:Npn \__regex_group_resetting:nnnN
25572 { \__regex_show_group_aux:nnnnN { ~(resetting) } }
25573 \cs_set_eq:NN \__regex_class:NnnnN \__regex_show_class:NnnnN
25574 \cs_set_protected:Npn \__regex_command_K:
25575 { \__regex_show_one:n { reset~match~start~(\iow_char:N\\K) } }
25576 \cs_set_protected:Npn \__regex_assertion:Nn ##1##2
25577 {
25578   \__regex_show_one:n
25579   { \bool_if:NF ##1 { negative~ } assertion:~##2 }
25580 }
25581 \cs_set:Npn \__regex_b_test: { word~boundary }
25582 \cs_set:Npn \__regex_Z_test: { anchor~at~end~(\iow_char:N\\Z) }
25583 \cs_set:Npn \__regex_A_test: { anchor~at~start~(\iow_char:N\\A) }
25584 \cs_set:Npn \__regex_G_test: { anchor~at~start~of~match~(\iow_char:N\\G) }
25585 \cs_set_protected:Npn \__regex_item_caseful_equal:n ##1
25586 { \__regex_show_one:n { char~code~\int_eval:n{##1} } }
25587 \cs_set_protected:Npn \__regex_item_caseful_range:nn ##1##2
25588 {
25589   \__regex_show_one:n
25590   { range~[\int_eval:n{##1}, \int_eval:n{##2}] }
25591 }
25592 \cs_set_protected:Npn \__regex_item_caseless_equal:n ##1
25593 { \__regex_show_one:n { char~code~\int_eval:n{##1}~(caseless) } }
25594 \cs_set_protected:Npn \__regex_item_caseless_range:nn ##1##2
25595 {
25596   \__regex_show_one:n
25597   { Range~[\int_eval:n{##1}, \int_eval:n{##2}]~(caseless) }
25598 }
25599 \cs_set_protected:Npn \__regex_item_catcode:nT
25600 { \__regex_show_item_catcode:NnT \c_true_bool }
25601 \cs_set_protected:Npn \__regex_item_catcode_reverse:nT
25602 { \__regex_show_item_catcode:NnT \c_false_bool }
25603 \cs_set_protected:Npn \__regex_item_reverse:n
25604 { \__regex_show_scope:nn { Reversed~match } }
25605 \cs_set_protected:Npn \__regex_item_exact:nn ##1##2
25606 { \__regex_show_one:n { char~##2,~catcode~##1 } }
25607 \cs_set_eq:NN \__regex_item_exact_cs:n \__regex_show_item_exact_cs:n
25608 \cs_set_protected:Npn \__regex_item_cs:n
25609 { \__regex_show_scope:nn { control~sequence } }

```

```

25610     \cs_set:cpn { __regex_prop_.: } { \__regex_show_one:n { any~token } }
25611     \seq_clear:N \l__regex_show_prefix_seq
25612     \__regex_show_push:n { ~ }
25613     \cs_if_exist_use:N #1
25614     \tl_build_end:N \l__regex_build_tl
25615     \exp_args:NNNo
25616     \group_end:
25617     \tl_set:Nn \l__regex_internal_a_tl { \l__regex_build_tl }
25618 }

```

(End definition for __regex_show:N.)

__regex_show_one:n Every part of the final message go through this function, which adds one line to the output, with the appropriate prefix.

```

25619 \cs_new_protected:Npn \__regex_show_one:n #1
25620 {
25621     \int_incr:N \l__regex_show_lines_int
25622     \tl_build_put_right:Nx \l__regex_build_tl
25623     {
25624         \exp_not:N \iow_newline:
25625         \seq_map_function:NN \l__regex_show_prefix_seq \use:n
25626         #1
25627     }
25628 }

```

(End definition for __regex_show_one:n.)

__regex_show_push:n Enter and exit levels of nesting. The scope function prints its first argument as an “introduction”, then performs its second argument in a deeper level of nesting.

```

\__regex_show_pop:
\__regex_show_scope:nn
25629 \cs_new_protected:Npn \__regex_show_push:n #1
25630 { \seq_put_right:Nx \l__regex_show_prefix_seq { #1 ~ } }
25631 \cs_new_protected:Npn \__regex_show_pop:
25632 { \seq_pop_right:NN \l__regex_show_prefix_seq \l__regex_internal_a_tl }
25633 \cs_new_protected:Npn \__regex_show_scope:nn #1#2
25634 {
25635     \__regex_show_one:n {#1}
25636     \__regex_show_push:n { ~ }
25637     #2
25638     \__regex_show_pop:
25639 }

```

(End definition for __regex_show_push:n, __regex_show_pop:, and __regex_show_scope:nn.)

__regex_show_group_aux:nnnnN We display all groups in the same way, simply adding a message, (no capture) or (resetting), to special groups. The odd \use_ii:nn avoids printing a spurious +-branch for the first branch.

```

25640 \cs_new_protected:Npn \__regex_show_group_aux:nnnnN #1#2#3#4#5
25641 {
25642     \__regex_show_one:n { ,~group~begin #1 }
25643     \__regex_show_push:n { | }
25644     \use_ii:nn #2
25645     \__regex_show_pop:
25646     \__regex_show_one:n
25647     { ‘~group~end \__regex_msg_repeated:nnN {#3} {#4} #5 }
25648 }

```


(End definition for _regex_show_group_aux:nnnnN.)

_regex_show_class:NnnnN

I'm entirely unhappy about this function: I couldn't find a way to test if a class is a single test. Instead, collect the representation of the tests in the class. If that had more than one line, write Match or Don't match on its own line, with the repeating information if any. Then the various tests on lines of their own, and finally a line. Otherwise, we need to evaluate the representation of the tests again (since the prefix is incorrect). That's clunky, but not too expensive, since it's only one test.

```

25649 \cs_set:Npn \_regex_show_class:NnnnN #1#2#3#4#5
25650 {
25651   \group_begin:
25652     \tl_build_begin:N \l__regex_build_tl
25653     \int_zero:N \l__regex_show_lines_int
25654     \__regex_show_push:n {~}
25655     #2
25656     \int_compare:nTF { \l__regex_show_lines_int = 0 }
25657     {
25658       \group_end:
25659       \__regex_show_one:n { \bool_if:NTF #1 { Fail } { Pass } }
25660     }
25661     {
25662       \bool_if:nTF
25663       { #1 && \int_compare_p:n { \l__regex_show_lines_int = 1 } }
25664       {
25665         \group_end:
25666         #2
25667         \tl_build_put_right:Nn \l__regex_build_tl
25668         { \__regex_msg_repeated:nnN {#3} {#4} #5 }
25669       }
25670       {
25671         \tl_build_end:N \l__regex_build_tl
25672         \exp_args:NNNo
25673         \group_end:
25674         \tl_set:Nn \l__regex_internal_a_tl \l__regex_build_tl
25675         \__regex_show_one:n
25676         {
25677           \bool_if:NTF #1 { Match } { Don't-match }
25678           \__regex_msg_repeated:nnN {#3} {#4} #5
25679         }
25680         \tl_build_put_right:Nx \l__regex_build_tl
25681         { \exp_not:o \l__regex_internal_a_tl }
25682       }
25683     }
25684   }

```

(End definition for _regex_show_class:NnnnnN.)

_regex_show_item_catcode:NnT

Produce a sequence of categories which the catcode bitmap #2 contains, and show it, indenting the tests on which this catcode constraint applies.

```

25685 \cs_new_protected:Npn \_regex_show_item_catcode:NnT #1#2
25686 {
25687   \seq_set_split:Nnn \l__regex_internal_seq { } { CBEMTPUDSLOA }
25688   \seq_set_filter:NNn \l__regex_internal_seq \l__regex_internal_seq
25689   { \int_if_odd_p:n { #2 / \int_use:c { c__regex_catcode_##1_int } } }

```

```

25690     \__regex_show_scope:nn
25691     {
25692         categories~
25693         \seq_map_function:NN \l__regex_internal_seq \use:n
25694         , ~
25695         \bool_if:NF #1 { negative~ } class
25696     }
25697 }

```

(End definition for __regex_show_item_catcode:NnT.)

__regex_show_item_exact_cs:n

```

25698 \cs_new_protected:Npn \__regex_show_item_exact_cs:n #1
25699 {
25700     \seq_set_split:Nnn \l__regex_internal_seq { \scan_stop: } {#1}
25701     \seq_set_map_x:Nnn \l__regex_internal_seq
25702         \l__regex_internal_seq { \iow_char:N\##1 }
25703     \__regex_show_one:n
25704     { control~sequence~ \seq_use:Nn \l__regex_internal_seq { ~or~ } }
25705 }

```

(End definition for __regex_show_item_exact_cs:n.)

41.4 Building

41.4.1 Variables used while building

\l__regex_min_state_int The last state that was allocated is \l__regex_max_state_int − 1, so that \l__regex_max_state_int always points to a free state. The min_state variable is 1 to begin with, but gets shifted in nested calls to the matching code, namely in \c{...} constructions.

```

25706 \int_new:N \l__regex_min_state_int
25707 \int_set:Nn \l__regex_min_state_int { 1 }
25708 \int_new:N \l__regex_max_state_int

```

(End definition for \l__regex_min_state_int and \l__regex_max_state_int.)

\l__regex_left_state_int Alternatives are implemented by branching from a left state into the various choices, then merging those into a right state. We store information about those states in two sequences. Those states are also used to implement group quantifiers. Most often, the left and right pointers only differ by 1.

```

25709 \int_new:N \l__regex_left_state_int
25710 \int_new:N \l__regex_right_state_int
25711 \seq_new:N \l__regex_left_state_seq
25712 \seq_new:N \l__regex_right_state_seq

```

(End definition for \l__regex_left_state_int and others.)

\l__regex_capturing_group_int

\l__regex_capturing_group_int is the next ID number to be assigned to a capturing group. This starts at 0 for the group enclosing the full regular expression, and groups are counted in the order of their left parenthesis, except when encountering resetting groups.

```

25713 \int_new:N \l__regex_capturing_group_int

```

(End definition for \l__regex_capturing_group_int.)

41.4.2 Framework

This phase is about going from a compiled regex to an NFA. Each state of the NFA is stored in a `\toks`. The operations which can appear in the `\toks` are

- `__regex_action_start_wildcard:N` $\langle boolean \rangle$ inserted at the start of the regular expression, where a `true` $\langle boolean \rangle$ makes it unanchored.
- `__regex_action_success:` marks the exit state of the NFA.
- `__regex_action_cost:n` $\{\langle shift \rangle\}$ is a transition from the current $\langle state \rangle$ to $\langle state \rangle + \langle shift \rangle$, which consumes the current character: the target state is saved and will be considered again when matching at the next position.
- `__regex_action_free:n` $\{\langle shift \rangle\}$, and `__regex_action_free_group:n` $\{\langle shift \rangle\}$ are free transitions, which immediately perform the actions for the state $\langle state \rangle + \langle shift \rangle$ of the NFA. They differ in how they detect and avoid infinite loops. For now, we just need to know that the `group` variant must be used for transitions back to the start of a group.
- `__regex_action_submatch:nN` $\{\langle group \rangle\}$ $\langle key \rangle$ where the $\langle key \rangle$ is `<` or `>` for the beginning or end of group numbered $\langle group \rangle$. This causes the current position in the query to be stored as the $\langle key \rangle$ submatch boundary.
- One of these actions, within a conditional.

We strive to preserve the following properties while building.

- The current capturing group is `capturing_group - 1`, and if a group opened now it would be labelled `capturing_group`.
- The last allocated state is `max_state - 1`, so `max_state` is a free state.
- The `left_state` points to a state to the left of the current group or of the last class.
- The `right_state` points to a newly created, empty state, with some transitions leading to it.
- The `left/right` sequences hold a list of the corresponding end-points of nested groups.

<code>__regex_build:n</code> <code>__regex_build_aux:Nn</code> <code>__regex_build:N</code> <code>__regex_build_aux:NN</code>	<p>The <code>n</code>-type function first compiles its argument. Reset some variables. Allocate two states, and put a wildcard in state 0 (transitions to state 1 and 0 state). Then build the regex within a (capturing) group numbered 0 (current value of <code>capturing_group</code>). Finally, if the match reaches the last state, it is successful. A <code>false</code> boolean for argument <code>#1</code> for the auxiliaries will suppress the wildcard and make the match anchored: used for <code>\peek_regex:nTF</code> and similar.</p>
--	--

```

25714 \cs_new_protected:Npn \__regex_build:n
25715 { \__regex_build_aux:Nn \c_true_bool }
25716 \cs_new_protected:Npn \__regex_build:N
25717 { \__regex_build_aux:NN \c_true_bool }
25718 \cs_new_protected:Npn \__regex_build_aux:Nn #1#2
25719 {
25720   \__regex_compile:n {#2}

```

```

25721   \_regex_build_aux:NN #1 \l__regex_internal_regex
25722   }
25723 \cs_new_protected:Npn \_regex_build_aux:NN #1#2
25724   {
25725     \_regex_standard_escapechar:
25726     \int_zero:N \l__regex_capturing_group_int
25727     \int_set_eq:NN \l__regex_max_state_int \l__regex_min_state_int
25728     \_regex_build_new_state:
25729     \_regex_build_new_state:
25730     \_regex_toks_put_right:Nn \l__regex_left_state_int
25731     { \_regex_action_start_wildcard:N #1 }
25732     \_regex_group:nnnN {#2} { 1 } { 0 } \c_false_bool
25733     \_regex_toks_put_right:Nn \l__regex_right_state_int
25734     { \_regex_action_success: }
25735   }

```

(End definition for `_regex_build:n` and others.)

`_regex_build_for_cs:n` The matching code relies on some global intarray variables, but only uses a range of their entries. Specifically,

- `\g__regex_state_active_intarray` from `\l__regex_min_state_int` to `\l__regex_max_state_int`;

Here, in this nested call to the matching code, we need the new versions of this range to involve completely new entries of the intarray variables, so we begin by setting (the new) `\l__regex_min_state_int` to (the old) `\l__regex_max_state_int` to use higher entries.

When using a regex to match a cs, we don't insert a wildcard, we anchor at the end, and since we ignore submatches, there is no need to surround the expression with a group. However, for branches to work properly at the outer level, we need to put the appropriate `left` and `right` states in their sequence.

```

25736 \cs_new_protected:Npn \_regex_build_for_cs:n #1
25737   {
25738     \int_set_eq:NN \l__regex_min_state_int \l__regex_max_state_int
25739     \_regex_build_new_state:
25740     \_regex_build_new_state:
25741     \_regex_push_lr_states:
25742     #1
25743     \_regex_pop_lr_states:
25744     \_regex_toks_put_right:Nn \l__regex_right_state_int
25745     {
25746       \if_int_compare:w -2 = \l__regex_curr_char_int
25747       \exp_after:wN \_regex_action_success:
25748       \fi:
25749     }
25750   }

```

(End definition for `_regex_build_for_cs:n`.)

41.4.3 Helpers for building an nfa

`_regex_push_lr_states:` When building the regular expression, we keep track of pointers to the left-end and right-end of each group without help from TEX's grouping.

```

25751 \cs_new_protected:Npn \__regex_push_lr_states:
25752 {
25753   \seq_push:No \l__regex_left_state_seq
25754   { \int_use:N \l__regex_left_state_int }
25755   \seq_push:No \l__regex_right_state_seq
25756   { \int_use:N \l__regex_right_state_int }
25757 }
25758 \cs_new_protected:Npn \__regex_pop_lr_states:
25759 {
25760   \seq_pop:NN \l__regex_left_state_seq \l__regex_internal_a_tl
25761   \int_set:Nn \l__regex_left_state_int \l__regex_internal_a_tl
25762   \seq_pop:NN \l__regex_right_state_seq \l__regex_internal_a_tl
25763   \int_set:Nn \l__regex_right_state_int \l__regex_internal_a_tl
25764 }

```

(End definition for __regex_push_lr_states: and __regex_pop_lr_states:.)

__regex_build_transition_left:NNN
 __regex_build_transition_right:nNn

Add a transition from #2 to #3 using the function #1. The `left` function is used for higher priority transitions, and the `right` function for lower priority transitions (which should be performed later). The signatures differ to reflect the differing usage later on. Both functions could be optimized.

```

25765 \cs_new_protected:Npn \__regex_build_transition_left:NNN #1#2#3
25766 { \__regex_toks_put_left:Nx #2 { #1 { \int_eval:n { #3 - #2 } } } }
25767 \cs_new_protected:Npn \__regex_build_transition_right:nNn #1#2#3
25768 { \__regex_toks_put_right:Nx #2 { #1 { \int_eval:n { #3 - #2 } } } }

```

(End definition for __regex_build_transition_left:NNN and __regex_build_transition_right:nNn.)

__regex_build_new_state:

Add a new empty state to the NFA. Then update the `left`, `right`, and `max` states, so that the `right` state is the new empty state, and the `left` state points to the previously “current” state.

```

25769 \cs_new_protected:Npn \__regex_build_new_state:
25770 {
25771   \__regex_toks_clear:N \l__regex_max_state_int
25772   \int_set_eq:NN \l__regex_left_state_int \l__regex_right_state_int
25773   \int_set_eq:NN \l__regex_right_state_int \l__regex_max_state_int
25774   \int_incr:N \l__regex_max_state_int
25775 }

```

(End definition for __regex_build_new_state:.)

__regex_build_transitions_lazyness:NNNNN

This function creates a new state, and puts two transitions starting from the old current state. The order of the transitions is controlled by #1, true for lazy quantifiers, and false for greedy quantifiers.

```

25776 \cs_new_protected:Npn \__regex_build_transitions_lazyness:NNNNN #1#2#3#4#5
25777 {
25778   \__regex_build_new_state:
25779   \__regex_toks_put_right:Nx \l__regex_left_state_int
25780   {
25781     \if_meaning:w \c_true_bool #1
25782     #2 { \int_eval:n { #3 - \l__regex_left_state_int } }
25783     #4 { \int_eval:n { #5 - \l__regex_left_state_int } }
25784     \else:
25785     #4 { \int_eval:n { #5 - \l__regex_left_state_int } }

```

```

25786         #2 { \int_eval:n { #3 - \l__regex_left_state_int } }
25787     \fi:
25788 }
25789 }

```

(End definition for `__regex_build_transitions_lazyness:NNNNN`.)

41.4.4 Building classes

`__regex_class:NnnnN`
`__regex_tests_action_cost:n`

The arguments are: $\langle\text{boolean}\rangle$ $\{\langle\text{tests}\rangle\}$ $\{\langle\text{min}\rangle\}$ $\{\langle\text{more}\rangle\}$ $\langle\text{lazyness}\rangle$. First store the tests with a trailing `__regex_action_cost:n`, in the true branch of `__regex_break_point:TF` for positive classes, or the false branch for negative classes. The integer $\langle\text{more}\rangle$ is 0 for fixed repetitions, -1 for unbounded repetitions, and $\langle\text{max}\rangle - \langle\text{min}\rangle$ for a range of repetitions.

```

25790 \cs_new_protected:Npn \__regex_class:NnnnN #1#2#3#4#5
25791 {
25792     \cs_set:Npx \__regex_tests_action_cost:n ##1
25793     {
25794         \exp_not:n { \exp_not:n {#2} }
25795         \bool_if:NTF #1
25796             { \__regex_break_point:TF { \__regex_action_cost:n {##1} } { } }
25797             { \__regex_break_point:TF { } { \__regex_action_cost:n {##1} } }
25798     }
25799     \if_case:w - #4 \exp_stop_f:
25800         \__regex_class_repeat:n {#3}
25801     \or: \__regex_class_repeat:nN {#3} #5
25802     \else: \__regex_class_repeat:nnN {#3} {#4} #5
25803     \fi:
25804 }
25805 \cs_new:Npn \__regex_tests_action_cost:n { \__regex_action_cost:n }

```

(End definition for `__regex_class:NnnnN` and `__regex_tests_action_cost:n`.)

`__regex_class_repeat:n`

This is used for a fixed number of repetitions. Build one state for each repetition, with a transition controlled by the tests that we have collected. That works just fine for $\#1 = 0$ repetitions: nothing is built.

```

25806 \cs_new_protected:Npn \__regex_class_repeat:n #1
25807 {
25808     \prg_replicate:nn {#1}
25809     {
25810         \__regex_build_new_state:
25811         \__regex_build_transition_right:nNn \__regex_tests_action_cost:n
25812         \l__regex_left_state_int \l__regex_right_state_int
25813     }
25814 }

```

(End definition for `__regex_class_repeat:n`.)

`__regex_class_repeat:nN`

This implements unbounded repetitions of a single class (e.g. the $*$ and $+$ quantifiers). If the minimum number $\#1$ of repetitions is 0, then build a transition from the current state to itself governed by the tests, and a free transition to a new state (hence skipping the tests). Otherwise, call `__regex_class_repeat:n` for the code to match $\#1$ repetitions, and add free transitions from the last state to the previous one, and to a new one. In both cases, the order of transitions is controlled by the lazyness boolean $\#2$.

```

25815 \cs_new_protected:Npn \__regex_class_repeat:nn #1#2
25816 {
25817   \if_int_compare:w #1 = 0 \exp_stop_f:
25818     \__regex_build_transitions_lazyness:NNNN #2
25819     \__regex_action_free:n \l__regex_right_state_int
25820     \__regex_tests_action_cost:n \l__regex_left_state_int
25821   \else:
25822     \__regex_class_repeat:n {#1}
25823     \int_set_eq:NN \l__regex_internal_a_int \l__regex_left_state_int
25824     \__regex_build_transitions_lazyness:NNNN #2
25825     \__regex_action_free:n \l__regex_right_state_int
25826     \__regex_action_free:n \l__regex_internal_a_int
25827   \fi:
25828 }

```

(End definition for __regex_class_repeat:nn.)

__regex_class_repeat:nnN We want to build the code to match from #1 to #1 + #2 repetitions. Match #1 repetitions (can be 0). Compute the final state of the next construction as a. Build #2 > 0 states, each with a transition to the next state governed by the tests, and a transition to the final state a. The computation of a is safe because states are allocated in order, starting from max_state.

```

25829 \cs_new_protected:Npn \__regex_class_repeat:nnN #1#2#3
25830 {
25831   \__regex_class_repeat:n {#1}
25832   \int_set:Nn \l__regex_internal_a_int
25833     { \l__regex_max_state_int + #2 - 1 }
25834   \prg_replicate:nn { #2 }
25835   {
25836     \__regex_build_transitions_lazyness:NNNN #3
25837     \__regex_action_free:n \l__regex_internal_a_int
25838     \__regex_tests_action_cost:n \l__regex_right_state_int
25839   }
25840 }

```

(End definition for __regex_class_repeat:nnN.)

41.4.5 Building groups

__regex_group_aux:nnnnN Arguments: {<label>} {<contents>} {<min>} {<more>} <lazyness>. If <min> is 0, we need to add a state before building the group, so that the thread which skips the group does not also set the start-point of the submatch. After adding one more state, the left_state is the left end of the group, from which all branches stem, and the right_state is the right end of the group, and all branches end their course in that state. We store those two integers to be queried for each branch, we build the NFA states for the contents #2 of the group, and we forget about the two integers. Once this is done, perform the repetition: either exactly #3 times, or #3 or more times, or between #3 and #3 + #4 times, with lazyness #5. The <label> #1 is used for submatch tracking. Each of the three auxiliaries expects left_state and right_state to be set properly.

```

25841 \cs_new_protected:Npn \__regex_group_aux:nnnnN #1#2#3#4#5
25842 {
25843   \if_int_compare:w #3 = 0 \exp_stop_f:
25844     \__regex_build_new_state:

```

```

25845 (assert)\assert_int:n { \l__regex_max_state_int = \l__regex_right_state_int + 1 }
25846     \__regex_build_transition_right:nNn \__regex_action_free_group:n
25847     \l__regex_left_state_int \l__regex_right_state_int
25848     \fi:
25849     \__regex_build_new_state:
25850     \__regex_push_lr_states:
25851     #2
25852     \__regex_pop_lr_states:
25853     \if_case:w - #4 \exp_stop_f:
25854         \__regex_group_repeat:nn {#1} {#3}
25855     \or: \__regex_group_repeat:nnN {#1} {#3} #5
25856     \else: \__regex_group_repeat:nnnN {#1} {#3} {#4} #5
25857     \fi:
25858 }

```

(End definition for __regex_group_aux:nnnnN.)

__regex_group:nnnN Hand to __regex_group_aux:nnnnN the label of that group (expanded), and the group itself, with some extra commands to perform.

__regex_group_no_capture:nnnN

```

25859 \cs_new_protected:Npn \__regex_group:nnnN #1
25860 {
25861     \exp_args:No \__regex_group_aux:nnnnN
25862     { \int_use:N \l__regex_capturing_group_int }
25863     {
25864         \int_incr:N \l__regex_capturing_group_int
25865         #1
25866     }
25867 }
25868 \cs_new_protected:Npn \__regex_group_no_capture:nnnN
25869 { \__regex_group_aux:nnnnN { -1 } }

```

(End definition for __regex_group:nnnN and __regex_group_no_capture:nnnN.)

__regex_group_resetting:nnnN Again, hand the label -1 to __regex_group_aux:nnnnN, but this time we work a little bit harder to keep track of the maximum group label at the end of any branch, and to reset the group number at each branch. This relies on the fact that a compiled regex always is a sequence of items of the form __regex_branch:n {<branch>}.

__regex_group_resetting_loop:nnNn

```

25870 \cs_new_protected:Npn \__regex_group_resetting:nnnN #1
25871 {
25872     \__regex_group_aux:nnnnN { -1 }
25873     {
25874         \exp_args:Noo \__regex_group_resetting_loop:nnNn
25875         { \int_use:N \l__regex_capturing_group_int }
25876         { \int_use:N \l__regex_capturing_group_int }
25877         #1
25878         { ?? \prg_break:n } { }
25879         \prg_break_point:
25880     }
25881 }
25882 \cs_new_protected:Npn \__regex_group_resetting_loop:nnNn #1#2#3#4
25883 {
25884     \use_none:nn #3 { \int_set:Nn \l__regex_capturing_group_int {#1} }
25885     \int_set:Nn \l__regex_capturing_group_int {#2}
25886     #3 {#4}

```



```

25887 \exp_args:Nf \__regex_group_resetting_loop:nnNn
25888 { \int_max:nn {#1} { \l__regex_capturing_group_int } }
25889 {#2}
25890 }

```

(End definition for __regex_group_resetting:nnnN and __regex_group_resetting_loop:nnNn.)

`__regex_branch:n` Add a free transition from the left state of the current group to a brand new state, starting point of this branch. Once the branch is built, add a transition from its last state to the right state of the group. The left and right states of the group are extracted from the relevant sequences.

```

25891 \cs_new_protected:Npn \__regex_branch:n #1
25892 {
25893   \__regex_build_new_state:
25894   \seq_get:NN \l__regex_left_state_seq \l__regex_internal_a_tl
25895   \int_set:Nn \l__regex_left_state_int \l__regex_internal_a_tl
25896   \__regex_build_transition_right:nNn \__regex_action_free:n
25897   \l__regex_left_state_int \l__regex_right_state_int
25898   #1
25899   \seq_get:NN \l__regex_right_state_seq \l__regex_internal_a_tl
25900   \__regex_build_transition_right:nNn \__regex_action_free:n
25901   \l__regex_right_state_int \l__regex_internal_a_tl
25902 }

```

(End definition for __regex_branch:n.)

`__regex_group_repeat:nn` This function is called to repeat a group a fixed number of times #2; if this is 0 we remove the group altogether (but don't reset the `capturing_group` label). Otherwise, the auxiliary `__regex_group_repeat_aux:n` copies #2 times the `\toks` for the group, and leaves `internal_a` pointing to the left end of the last repetition. We only record the submatch information at the last repetition. Finally, add a state at the end (the transition to it has been taken care of by the replicating auxiliary).

```

25903 \cs_new_protected:Npn \__regex_group_repeat:nn #1#2
25904 {
25905   \if_int_compare:w #2 = 0 \exp_stop_f:
25906   \int_set:Nn \l__regex_max_state_int
25907   { \l__regex_left_state_int - 1 }
25908   \__regex_build_new_state:
25909   \else:
25910   \__regex_group_repeat_aux:n {#2}
25911   \__regex_group_submatches:nnN {#1}
25912   \l__regex_internal_a_int \l__regex_right_state_int
25913   \__regex_build_new_state:
25914   \fi:
25915 }

```

(End definition for __regex_group_repeat:nn.)

`__regex_group_submatches:nnN` This inserts in states #2 and #3 the code for tracking submatches of the group #1, unless inhibited by a label of -1.

```

25916 \cs_new_protected:Npn \__regex_group_submatches:nnN #1#2#3
25917 {
25918   \if_int_compare:w #1 > - 1 \exp_stop_f:
25919   \__regex_toks_put_left:Nx #2 { \__regex_action_submatch:nN {#1} < }

```

```

25920     \__regex_toks_put_left:Nx #3 { \__regex_action_submatch:nN {#1} > }
25921     \fi:
25922 }

```

(End definition for __regex_group_submatches:nNN.)

__regex_group_repeat_aux:n Here we repeat \toks ranging from left_state to max_state, #1 > 0 times. First add a transition so that the copies “chain” properly. Compute the shift c between the original copy and the last copy we want. Shift the right_state and max_state to their final values. We then want to perform c copy operations. At the end, b is equal to the max_state, and a points to the left of the last copy of the group.

```

25923 \cs_new_protected:Npn \__regex_group_repeat_aux:n #1
25924 {
25925     \__regex_build_transition_right:nNn \__regex_action_free:n
25926     \l__regex_right_state_int \l__regex_max_state_int
25927     \int_set_eq:NN \l__regex_internal_a_int \l__regex_left_state_int
25928     \int_set_eq:NN \l__regex_internal_b_int \l__regex_max_state_int
25929     \if_int_compare:w \int_eval:n {#1} > 1 \exp_stop_f:
25930     \int_set:Nn \l__regex_internal_c_int
25931     {
25932         ( #1 - 1 )
25933         * ( \l__regex_internal_b_int - \l__regex_internal_a_int )
25934     }
25935     \int_add:Nn \l__regex_right_state_int { \l__regex_internal_c_int }
25936     \int_add:Nn \l__regex_max_state_int { \l__regex_internal_c_int }
25937     \__regex_toks_memcpy:Nn
25938     \l__regex_internal_b_int
25939     \l__regex_internal_a_int
25940     \l__regex_internal_c_int
25941     \fi:
25942 }

```

(End definition for __regex_group_repeat_aux:n.)

__regex_group_repeat:nnN This function is called to repeat a group at least n times; the case n = 0 is very different from n > 0. Assume first that n = 0. Insert submatch tracking information at the start and end of the group, add a free transition from the right end to the “true” left state a (remember: in this case we had added an extra state before the left state). This forms the loop, which we break away from by adding a free transition from a to a new state.

Now consider the case n > 0. Repeat the group n times, chaining various copies with a free transition. Add submatch tracking only to the last copy, then add a free transition from the right end back to the left end of the last copy, either before or after the transition to move on towards the rest of the NFA. This transition can end up before submatch tracking, but that is irrelevant since it only does so when going again through the group, recording new matches. Finally, add a state; we already have a transition pointing to it from __regex_group_repeat_aux:n.

```

25943 \cs_new_protected:Npn \__regex_group_repeat:nnN #1#2#3
25944 {
25945     \if_int_compare:w #2 = 0 \exp_stop_f:
25946     \__regex_group_submatches:nNN {#1}
25947     \l__regex_left_state_int \l__regex_right_state_int
25948     \int_set:Nn \l__regex_internal_a_int
25949     { \l__regex_left_state_int - 1 }

```

```

25950     \__regex_build_transition_right:nNn \__regex_action_free:n
25951     \l__regex_right_state_int \l__regex_internal_a_int
25952     \__regex_build_new_state:
25953     \if_meaning:w \c_true_bool #3
25954     \__regex_build_transition_left:NNN \__regex_action_free:n
25955     \l__regex_internal_a_int \l__regex_right_state_int
25956     \else:
25957     \__regex_build_transition_right:nNn \__regex_action_free:n
25958     \l__regex_internal_a_int \l__regex_right_state_int
25959     \fi:
25960 \else:
25961     \__regex_group_repeat_aux:n {#2}
25962     \__regex_group_submatches:nNN {#1}
25963     \l__regex_internal_a_int \l__regex_right_state_int
25964     \if_meaning:w \c_true_bool #3
25965     \__regex_build_transition_right:nNn \__regex_action_free_group:n
25966     \l__regex_right_state_int \l__regex_internal_a_int
25967     \else:
25968     \__regex_build_transition_left:NNN \__regex_action_free_group:n
25969     \l__regex_right_state_int \l__regex_internal_a_int
25970     \fi:
25971     \__regex_build_new_state:
25972 \fi:
25973 }

```

(End definition for __regex_group_repeat:nnN.)

__regex_group_repeat:nnnN

We wish to repeat the group between #2 and #2 + #3 times, with a laziness controlled by #4. We insert submatch tracking up front: in principle, we could avoid recording submatches for the first #2 copies of the group, but that forces us to treat specially the case #2 = 0. Repeat that group with submatch tracking #2 + #3 times (the maximum number of repetitions). Then our goal is to add #3 transitions from the end of the #2-th group, and each subsequent groups, to the end. For a lazy quantifier, we add those transitions to the left states, before submatch tracking. For the greedy case, we add the transitions to the right states, after submatch tracking and the transitions which go on with more repetitions. In the greedy case with #2 = 0, the transition which skips over all copies of the group must be added separately, because its starting state does not follow the normal pattern: we had to add it “by hand” earlier.

```

25974 \cs_new_protected:Npn \__regex_group_repeat:nnnN #1#2#3#4
25975 {
25976     \__regex_group_submatches:nNN {#1}
25977     \l__regex_left_state_int \l__regex_right_state_int
25978     \__regex_group_repeat_aux:n { #2 + #3 }
25979     \if_meaning:w \c_true_bool #4
25980     \int_set_eq:NN \l__regex_left_state_int \l__regex_max_state_int
25981     \prg_replicate:nn { #3 }
25982     {
25983         \int_sub:Nn \l__regex_left_state_int
25984         { \l__regex_internal_b_int - \l__regex_internal_a_int }
25985         \__regex_build_transition_left:NNN \__regex_action_free:n
25986         \l__regex_left_state_int \l__regex_max_state_int
25987     }
25988 \else:

```

```

25989 \prg_replicate:nn { #3 - 1 }
25990 {
25991   \int_sub:Nn \l__regex_right_state_int
25992   { \l__regex_internal_b_int - \l__regex_internal_a_int }
25993   \__regex_build_transition_right:nNn \__regex_action_free:n
25994   \l__regex_right_state_int \l__regex_max_state_int
25995 }
25996 \if_int_compare:w #2 = 0 \exp_stop_f:
25997   \int_set:Nn \l__regex_right_state_int
25998   { \l__regex_left_state_int - 1 }
25999 \else:
26000   \int_sub:Nn \l__regex_right_state_int
26001   { \l__regex_internal_b_int - \l__regex_internal_a_int }
26002 \fi:
26003 \__regex_build_transition_right:nNn \__regex_action_free:n
26004 \l__regex_right_state_int \l__regex_max_state_int
26005 \fi:
26006 \__regex_build_new_state:
26007 }

```

(End definition for __regex_group_repeat:nnnN.)

41.4.6 Others

__regex_assertion:Nn Usage: __regex_assertion:Nn *<boolean>* {*<test>*}, where the *<test>* is either of the two other functions. Add a free transition to a new state, conditionally to the assertion test. The __regex_b_test: test is used by the \b and \B escape: check if the last character was a word character or not, and do the same to the current character. The __regex_A_test: test is used by the \A escape: check if the last boundary-markers of the string are non-word characters for this purpose. The __regex_G_test: test is used by the \G escape: check if the last character was a word character or not, and do the same to the current character. The __regex_Z_test: test is used by the \Z escape: check if the last character was a word character or not, and do the same to the current character.

```

26008 \cs_new_protected:Npn \__regex_assertion:Nn #1#2
26009 {
26010   \__regex_build_new_state:
26011   \__regex_toks_put_right:Nx \l__regex_left_state_int
26012   {
26013     \exp_not:n {#2}
26014     \__regex_break_point:TF
26015     \bool_if:NF #1 { { } }
26016     {
26017       \__regex_action_free:n
26018       {
26019         \int_eval:n
26020         { \l__regex_right_state_int - \l__regex_left_state_int }
26021       }
26022     }
26023     \bool_if:NT #1 { { } }
26024   }
26025 }
26026 \cs_new_protected:Npn \__regex_b_test:
26027 {
26028   \group_begin:
26029   \int_set_eq:NN \l__regex_curr_char_int \l__regex_last_char_int
26030   \__regex_prop_w:
26031   \__regex_break_point:TF
26032   { \group_end: \__regex_item_reverse:n \__regex_prop_w: }

```

```

26033         { \group_end: \_regex_prop_w: }
26034     }
26035 \cs_new_protected:Npn \_regex_Z_test:
26036 {
26037     \if_int_compare:w -2 = \l__regex_curr_char_int
26038         \exp_after:wN \_regex_break_true:w
26039     \fi:
26040 }
26041 \cs_new_protected:Npn \_regex_A_test:
26042 {
26043     \if_int_compare:w -2 = \l__regex_last_char_int
26044         \exp_after:wN \_regex_break_true:w
26045     \fi:
26046 }
26047 \cs_new_protected:Npn \_regex_G_test:
26048 {
26049     \if_int_compare:w \l__regex_curr_pos_int = \l__regex_start_pos_int
26050         \exp_after:wN \_regex_break_true:w
26051     \fi:
26052 }

```

(End definition for `_regex_assertion:Nn` and others.)

`_regex_command_K:` Change the starting point of the 0-th submatch (full match), and transition to a new state, pretending that this is a fresh thread.

```

26053 \cs_new_protected:Npn \_regex_command_K:
26054 {
26055     \_regex_build_new_state:
26056     \_regex_toks_put_right:Nx \l__regex_left_state_int
26057     {
26058         \_regex_action_submatch:nN { 0 } <
26059         \bool_set_true:N \l__regex_fresh_thread_bool
26060         \_regex_action_free:n
26061         {
26062             \int_eval:n
26063             { \l__regex_right_state_int - \l__regex_left_state_int }
26064         }
26065         \bool_set_false:N \l__regex_fresh_thread_bool
26066     }
26067 }

```

(End definition for `_regex_command_K:.`)

41.5 Matching

We search for matches by running all the execution threads through the NFA in parallel, reading one token of the query at each step. The NFA contains “free” transitions to other states, and transitions which “consume” the current token. For free transitions, the instruction at the new state of the NFA is performed immediately. When a transition consumes a character, the new state is appended to a list of “active states”, stored in `\g__regex_thread_info_intarray` (together with submatch information): this thread is made active again when the next token is read from the query. At every step (for each token in the query), we unpack that list of active states and the corresponding submatch props, and empty those.

If two paths through the NFA “collide” in the sense that they reach the same state after reading a given token, then they only differ in how they previously matched, and any future execution would be identical for both. (Note that this would be wrong in the presence of back-references.) Hence, we only need to keep one of the two threads: the thread with the highest priority. Our NFA is built in such a way that higher priority actions always come before lower priority actions, which makes things work.

The explanation in the previous paragraph may make us think that we simply need to keep track of which states were visited at a given step: after all, the loop generated when matching `(a?)*` against `a` is broken, isn’t it? No. The group first matches `a`, as it should, then repeats; it attempts to match `a` again but fails; it skips `a`, and finds out that this state has already been seen at this position in the query: the match stops. The capturing group is (wrongly) `a`. What went wrong is that a thread collided with itself, and the later version, which has gone through the group one more times with an empty match, should have a higher priority than not going through the group.

We solve this by distinguishing “normal” free transitions `__regex_action_free:n` from transitions `__regex_action_free_group:n` which go back to the start of the group. The former keeps threads unless they have been visited by a “completed” thread, while the latter kind of transition also prevents going back to a state visited by the current thread.

41.5.1 Variables used when matching

`\l__regex_min_pos_int` The tokens in the query are indexed from `min_pos` for the first to `max_pos - 1` for the last, and their information is stored in several arrays and `\toks` registers with those numbers.
`\l__regex_max_pos_int` We match without backtracking, keeping all threads in lockstep at the `curr_pos` in the query. The starting point of the current match attempt is `start_pos`, and `success_pos`,
`\l__regex_curr_pos_int` updated whenever a thread succeeds, is used as the next starting position.
`\l__regex_start_pos_int`
`\l__regex_success_pos_int`

```
26068 \int_new:N \l__regex_min_pos_int
26069 \int_new:N \l__regex_max_pos_int
26070 \int_new:N \l__regex_curr_pos_int
26071 \int_new:N \l__regex_start_pos_int
26072 \int_new:N \l__regex_success_pos_int
```

(End definition for `\l__regex_min_pos_int` and others.)

`\l__regex_curr_char_int` The character and category codes of the token at the current position and a token list
`\l__regex_curr_catcode_int` expanding to that token; the character code of the token at the previous position; the
`\l__regex_curr_token_tl` character code of the token just before a successful match; and the character code of the
`\l__regex_last_char_int` result of changing the case of the current token (`A-Z↔a-z`). This last integer is only
`\l__regex_last_char_success_int` computed when necessary, and is otherwise `\c_max_int`. The `curr_char` variable is also
`\l__regex_case_changed_char_int` used in various other phases to hold a character code.

```
26073 \int_new:N \l__regex_curr_char_int
26074 \int_new:N \l__regex_curr_catcode_int
26075 \tl_new:N \l__regex_curr_token_tl
26076 \int_new:N \l__regex_last_char_int
26077 \int_new:N \l__regex_last_char_success_int
26078 \int_new:N \l__regex_case_changed_char_int
```

(End definition for `\l__regex_curr_char_int` and others.)

`\l__regex_curr_state_int` For every character in the token list, each of the active states is considered in turn. The variable `\l__regex_curr_state_int` holds the state of the NFA which is currently considered: transitions are then given as shifts relative to the current state.

```
26079 \int_new:N \l__regex_curr_state_int
```

(End definition for `\l__regex_curr_state_int`.)

`\l__regex_curr_submatches_tl`
`\l__regex_success_submatches_tl` The submatches for the thread which is currently active are stored in the `curr_submatches` list, which is almost a comma list, but ends with a comma. This list is stored by `__regex_store_state:n` into an intarray variable, to be retrieved when matching at the next position. When a thread succeeds, this list is copied to `\l__regex_success_submatches_tl`: only the last successful thread remains there.

```
26080 \tl_new:N \l__regex_curr_submatches_tl
```

```
26081 \tl_new:N \l__regex_success_submatches_tl
```

(End definition for `\l__regex_curr_submatches_tl` and `\l__regex_success_submatches_tl`.)

`\l__regex_step_int` This integer, always even, is increased every time a character in the query is read, and not reset when doing multiple matches. We store in `\g__regex_state_active_intarray` the last step in which each $\langle state \rangle$ in the NFA was encountered. This lets us break infinite loops by not visiting the same state twice in the same step. In fact, the step we store is equal to `step` when we have started performing the operations of `\toks\langle state \rangle`, but not finished yet. However, once we finish, we store `step + 1` in `\g__regex_state_active_intarray`. This is needed to track submatches properly (see building phase). The `step` is also used to attach each set of submatch information to a given iteration (and automatically discard it when it corresponds to a past step).

```
26082 \int_new:N \l__regex_step_int
```

(End definition for `\l__regex_step_int`.)

`\l__regex_min_thread_int`
`\l__regex_max_thread_int` All the currently active threads are kept in order of precedence in `\g__regex_thread_info_intarray` together with the corresponding submatch information. Data in this intarray is organized as blocks from `min_thread` (included) to `max_thread` (excluded). At the start of every step, the whole array is unpacked, so that the space can immediately be reused, and `max_thread` is reset to `min_thread`, effectively clearing the array.

```
26083 \int_new:N \l__regex_min_thread_int
```

```
26084 \int_new:N \l__regex_max_thread_int
```

(End definition for `\l__regex_min_thread_int` and `\l__regex_max_thread_int`.)

`\g__regex_state_active_intarray`
`\g__regex_thread_info_intarray` `\g__regex_state_active_intarray` stores the last $\langle step \rangle$ in which each $\langle state \rangle$ was active. `\g__regex_thread_info_intarray` stores threads to be considered in the next step, more precisely the states in which these threads are.

```
26085 \intarray_new:Nn \g__regex_state_active_intarray { 65536 }
```

```
26086 \intarray_new:Nn \g__regex_thread_info_intarray { 65536 }
```

(End definition for `\g__regex_state_active_intarray` and `\g__regex_thread_info_intarray`.)

`\l_regex_matched_analysis_tl`
`\l__regex_curr_analysis_tl` The list `\l__regex_curr_analysis_tl` consists of a brace group containing three brace groups corresponding to the current token, with the same syntax as `\tl_analysis_map_inline:nn`. The list `\l__regex_matched_analysis_tl` (constructed under the `tl_build` machinery) has one item for each token that has already been treated so far in a given match attempt: each item consists of three brace groups with the same syntax as `\tl_analysis_map_inline:nn`.

```
26087 \tl_new:N \l__regex_matched_analysis_tl
26088 \tl_new:N \l__regex_curr_analysis_tl
```

(End definition for `\l__regex_matched_analysis_tl` and `\l__regex_curr_analysis_tl`.)

`\l__regex_every_match_tl` Every time a match is found, this token list is used. For single matching, the token list is empty. For multiple matching, the token list is set to repeat the matching, after performing some operation which depends on the user function. See `__regex_single_match:` and `__regex_multi_match:n`.

```
26089 \tl_new:N \l__regex_every_match_tl
```

(End definition for `\l__regex_every_match_tl`.)

`\l_regex_fresh_thread_bool`
`\l_regex_empty_success_bool`
`_regex_if_two_empty_matches:F` When doing multiple matches, we need to avoid infinite loops where each iteration matches the same empty token list. When an empty token list is matched, the next successful match of the same empty token list is suppressed. We detect empty matches by setting `\l__regex_fresh_thread_bool` to `true` for threads which directly come from the start of the regex or from the `\K` command, and testing that boolean whenever a thread succeeds. The function `__regex_if_two_empty_matches:F` is redefined at every match attempt, depending on whether the previous match was empty or not: if it was, then the function must cancel a purported success if it is empty and at the same spot as the previous match; otherwise, we definitely don't have two identical empty matches, so the function is `\use:n`.

```
26090 \bool_new:N \l__regex_fresh_thread_bool
26091 \bool_new:N \l__regex_empty_success_bool
26092 \cs_new_eq:NN \_regex_if_two_empty_matches:F \use:n
```

(End definition for `\l__regex_fresh_thread_bool`, `\l__regex_empty_success_bool`, and `_regex_if_two_empty_matches:F`.)

`\g__regex_success_bool`
`\l__regex_saved_success_bool`
`\l__regex_match_success_bool` The boolean `\l__regex_match_success_bool` is true if the current match attempt was successful, and `\g__regex_success_bool` is true if there was at least one successful match. This is the only global variable in this whole module, but we would need it to be local when matching a control sequence with `\c{...}`. This is done by saving the global variable into `\l__regex_saved_success_bool`, which is local, hence not affected by the changes due to inner regex functions.

```
26093 \bool_new:N \g__regex_success_bool
26094 \bool_new:N \l__regex_saved_success_bool
26095 \bool_new:N \l__regex_match_success_bool
```

(End definition for `\g__regex_success_bool`, `\l__regex_saved_success_bool`, and `\l__regex_match_success_bool`.)

41.5.2 Matching: framework

`__regex_match:n` Initialize the variables that should be set once for each user function (even for multiple matches). Namely, the overall matching is not yet successful; none of the states should be marked as visited (`__regex_state_active_intarray`), and we start at step 0; we pretend that there was a previous match ending at the start of the query, which was not empty (to avoid smothering an empty match at the start). Once all this is set up, we are ready for the ride. Find the first match.

```

26096 \cs_new_protected:Npn \__regex_match:n #1
26097 {
26098   \__regex_match_init:
26099   \__regex_match_once_init:
26100   \tl_analysis_map_inline:nn {#1}
26101     { \__regex_match_one_token:nnN {##1} {##2} ##3 }
26102   \__regex_match_one_token:nnN { } { -2 } F
26103   \prg_break_point:Nn \__regex_maplike_break: { }
26104 }
26105 \cs_new_protected:Npn \__regex_match_cs:n #1
26106 {
26107   \int_set_eq:NN \l__regex_min_thread_int \l__regex_max_thread_int
26108   \__regex_match_init:
26109   \__regex_match_once_init:
26110   \str_map_inline:nn {#1}
26111     {
26112       \tl_if_blank:nTF {##1}
26113         { \__regex_match_one_token:nnN {##1} {'##1} A }
26114         { \__regex_match_one_token:nnN {##1} {'##1} C }
26115     }
26116   \__regex_match_one_token:nnN { } { -2 } F
26117   \prg_break_point:Nn \__regex_maplike_break: { }
26118 }
26119 \cs_new_protected:Npn \__regex_match_init:
26120 {
26121   \bool_gset_false:N \g__regex_success_bool
26122   \int_step_inline:nnn
26123     \l__regex_min_state_int { \l__regex_max_state_int - 1 }
26124     {
26125       \_kernel_intarray_gset:Nnn
26126       \g__regex_state_active_intarray {##1} { 1 }
26127     }
26128   \int_zero:N \l__regex_step_int
26129   \int_set:Nn \l__regex_min_pos_int { 2 }
26130   \int_set_eq:NN \l__regex_success_pos_int \l__regex_min_pos_int
26131   \int_set:Nn \l__regex_last_char_success_int { -2 }
26132   \tl_build_begin:N \l__regex_matched_analysis_tl
26133   \tl_clear:N \l__regex_curr_analysis_tl
26134   \int_set:Nn \l__regex_min_submatch_int { 1 }
26135   \int_set_eq:NN \l__regex_submatch_int \l__regex_min_submatch_int
26136   \bool_set_false:N \l__regex_empty_success_bool
26137 }

```

(End definition for `__regex_match:n`, `__regex_match_cs:n`, and `__regex_match_init:.`)

`__regex_match_once_init:` This function resets various variables used when finding one match. It is called before the

loop through characters, and every time we find a match, before searching for another match (this is controlled by the `every_match` token list).

First initialize some variables: set the conditional which detects identical empty matches; this match attempt starts at the previous `success_pos`, is not yet successful, and has no submatches yet; clear the array of active threads, and put the starting state 0 in it. We are then almost ready to read our first token in the query, but we actually start one position earlier than the start because `__regex_match_one_token:nnN` increments `__regex_curr_pos_int` and saves `__regex_curr_char_int` as the `last_char` so that word boundaries can be correctly identified.

```

26138 \cs_new_protected:Npn \__regex_match_once_init:
26139 {
26140     \if_meaning:w \c_true_bool \l__regex_empty_success_bool
26141     \cs_set:Npn \__regex_if_two_empty_matches:F
26142     {
26143         \int_compare:nNnF
26144             \l__regex_start_pos_int = \l__regex_curr_pos_int
26145     }
26146     \else:
26147         \cs_set_eq:NN \__regex_if_two_empty_matches:F \use:n
26148     \fi:
26149     \int_set_eq:NN \l__regex_start_pos_int \l__regex_success_pos_int
26150     \bool_set_false:N \l__regex_match_success_bool
26151     \tl_set:Nx \l__regex_curr_submatches_tl
26152         { \prg_replicate:nn { 2 * \l__regex_capturing_group_int } { 0 , } }
26153     \int_set_eq:NN \l__regex_max_thread_int \l__regex_min_thread_int
26154     \__regex_store_state:n { \l__regex_min_state_int }
26155     \int_set:Nn \l__regex_curr_pos_int
26156         { \l__regex_start_pos_int - 1 }
26157     \int_set_eq:NN \l__regex_curr_char_int \l__regex_last_char_success_int
26158     \tl_build_get:NN \l__regex_matched_analysis_tl \l__regex_internal_a_tl
26159     \exp_args:NNf \__regex_match_once_init_aux:
26160     \tl_map_inline:nn
26161         { \exp_after:wN \l__regex_internal_a_tl \l__regex_curr_analysis_tl }
26162         { \__regex_match_one_token:nnN ##1 }
26163     \prg_break_point:Nn \__regex_maplike_break: { }
26164 }
26165 \cs_new_protected:Npn \__regex_match_once_init_aux:
26166 {
26167     \tl_build_clear:N \l__regex_matched_analysis_tl
26168     \tl_clear:N \l__regex_curr_analysis_tl
26169 }

```

(End definition for `__regex_match_once_init:`)

`__regex_single_match:` For a single match, the overall success is determined by whether the only match attempt is a success. When doing multiple matches, the overall matching is successful as soon as any match succeeds. Perform the action #1, then find the next match.

```

26170 \cs_new_protected:Npn \__regex_single_match:
26171 {
26172     \tl_set:Nn \l__regex_every_match_tl
26173     {
26174         \bool_gset_eq:NN
26175         \g__regex_success_bool

```

```

26176         \l__regex_match_success_bool
26177         \__regex_maplike_break:
26178     }
26179 }
26180 \cs_new_protected:Npn \__regex_multi_match:n #1
26181 {
26182     \tl_set:Nn \l__regex_every_match_tl
26183     {
26184         \if_meaning:w \c_false_bool \l__regex_match_success_bool
26185             \exp_after:wN \__regex_maplike_break:
26186         \fi:
26187         \bool_gset_true:N \g__regex_success_bool
26188         #1
26189         \__regex_match_once_init:
26190     }
26191 }

```

(End definition for __regex_single_match: and __regex_multi_match:n.)

__regex_match_one_token:nnN At each new position, set some variables and get the new character and category from the query. Then unpack the array of active threads, and clear it by resetting its length (max_thread). This results in a sequence of __regex_use_state_and_submatches:w <state>,<submatch-clist>; and we consider those states one by one in order. As soon as a thread succeeds, exit the step, and, if there are threads to consider at the next position, and we have not reached the end of the string, repeat the loop. Otherwise, the last thread that succeeded is the match. We explain the fresh_thread business when describing __regex_action_wildcard:.

```

26192 \cs_new_protected:Npn \__regex_match_one_token:nnN #1#2#3
26193 {
26194     \int_add:Nn \l__regex_step_int { 2 }
26195     \int_incr:N \l__regex_curr_pos_int
26196     \int_set_eq:NN \l__regex_last_char_int \l__regex_curr_char_int
26197     \int_set_eq:NN \l__regex_case_changed_char_int \c_max_int
26198     \tl_set:Nn \l__regex_curr_token_tl {#1}
26199     \int_set:Nn \l__regex_curr_char_int {#2}
26200     \int_set:Nn \l__regex_curr_catcode_int { "#3 }
26201     \tl_build_put_right:Nx \l__regex_matched_analysis_tl
26202     { \exp_not:o \l__regex_curr_analysis_tl }
26203     \tl_set:Nn \l__regex_curr_analysis_tl { { {#1} {#2} #3 } }
26204     \use:x
26205     {
26206         \int_set_eq:NN \l__regex_max_thread_int \l__regex_min_thread_int
26207         \int_step_function:nnN
26208             { \l__regex_min_thread_int }
26209             { \l__regex_max_thread_int - 1 }
26210         \__regex_match_one_active:n
26211     }
26212     \prg_break_point:
26213     \bool_set_false:N \l__regex_fresh_thread_bool
26214     \if_int_compare:w \l__regex_max_thread_int > \l__regex_min_thread_int
26215         \if_int_compare:w -2 < \l__regex_curr_char_int
26216             \exp_after:wN \exp_after:wN \exp_after:wN \use_none:n
26217         \fi:
26218     \fi:

```

```

26219     \l__regex_every_match_t1
26220 }
26221 \cs_new:Npn \__regex_match_one_active:n #1
26222 {
26223     \__regex_use_state_and_submatches:w
26224     \__kernel_intarray_range_to_clist:Nnn
26225     \g__regex_thread_info_intarray
26226     { 1 + #1 * (\l__regex_capturing_group_int * 2 + 1) }
26227     { (1 + #1) * (\l__regex_capturing_group_int * 2 + 1) }
26228 ;
26229 }

```

(End definition for __regex_match_one_token:nnN and __regex_match_one_active:n.)

41.5.3 Using states of the nfa

__regex_use_state: Use the current NFA instruction. The state is initially marked as belonging to the current **step**: this allows normal free transition to repeat, but group-repeating transitions won't. Once we are done exploring all the branches it spawned, the state is marked as **step + 1**: any thread hitting it at that point will be terminated.

```

26230 \cs_new_protected:Npn \__regex_use_state:
26231 {
26232     \__kernel_intarray_gset:Nnn \g__regex_state_active_intarray
26233     { \l__regex_curr_state_int } { \l__regex_step_int }
26234     \__regex_toks_use:w \l__regex_curr_state_int
26235     \__kernel_intarray_gset:Nnn \g__regex_state_active_intarray
26236     { \l__regex_curr_state_int }
26237     { \int_eval:n { \l__regex_step_int + 1 } }
26238 }

```

(End definition for __regex_use_state:.)

__regex_use_state_and_submatches:w This function is called as one item in the array of active threads after that array has been unpacked for a new step. Update the `curr_state` and `curr_submatches` and use the state if it has not yet been encountered at this step.

```

26239 \cs_new_protected:Npn \__regex_use_state_and_submatches:w #1 , #2 ;
26240 {
26241     \int_set:Nn \l__regex_curr_state_int {#1}
26242     \if_int_compare:w
26243         \__kernel_intarray_item:Nn \g__regex_state_active_intarray
26244         { \l__regex_curr_state_int }
26245         < \l__regex_step_int
26246     \tl_set:Nn \l__regex_curr_submatches_tl { #2 , }
26247     \exp_after:wN \__regex_use_state:
26248     \fi:
26249     \scan_stop:
26250 }

```

(End definition for __regex_use_state_and_submatches:w.)

41.5.4 Actions when matching

`__regex_action_start_wildcard:N` For an unanchored match, state 0 has a free transition to the next and a costly one to itself, to repeat at the next position. To catch repeated identical empty matches, we need to know if a successful thread corresponds to an empty match. The instruction resetting `\l__regex_fresh_thread_bool` may be skipped by a successful thread, hence we had to add it to `__regex_match_one_token:nnN` too.

```

26251 \cs_new_protected:Npn \__regex_action_start_wildcard:N #1
26252 {
26253   \bool_set_true:N \l__regex_fresh_thread_bool
26254   \__regex_action_free:n {1}
26255   \bool_set_false:N \l__regex_fresh_thread_bool
26256   \bool_if:NT #1 { \__regex_action_cost:n {0} }
26257 }
```

(End definition for `__regex_action_start_wildcard:N`.)

`__regex_action_free:n` These functions copy a thread after checking that the NFA state has not already been used at this position. If not, store submatches in the new state, and insert the instructions for that state in the input stream. Then restore the old value of `\l__regex_curr_state_int` and of the current submatches. The two types of free transitions differ by how they test that the state has not been encountered yet: the `group` version is stricter, and will not use a state if it was used earlier in the current thread, hence forcefully breaking the loop, while the “normal” version will revisit a state even within the thread itself.

```

26258 \cs_new_protected:Npn \__regex_action_free:n
26259 { \__regex_action_free_aux:nn { > \l__regex_step_int \else: } }
26260 \cs_new_protected:Npn \__regex_action_free_group:n
26261 { \__regex_action_free_aux:nn { < \l__regex_step_int } }
26262 \cs_new_protected:Npn \__regex_action_free_aux:nn #1#2
26263 {
26264   \use:x
26265   {
26266     \int_add:Nn \l__regex_curr_state_int {#2}
26267     \exp_not:n
26268     {
26269       \if_int_compare:w
26270         \__kernel_intarray_item:Nn \g__regex_state_active_intarray
26271         { \l__regex_curr_state_int }
26272         #1
26273         \exp_after:wN \__regex_use_state:
26274         \fi:
26275     }
26276     \int_set:Nn \l__regex_curr_state_int
26277     { \int_use:N \l__regex_curr_state_int }
26278     \tl_set:Nn \exp_not:N \l__regex_curr_submatches_tl
26279     { \exp_not:o \l__regex_curr_submatches_tl }
26280   }
26281 }
```

(End definition for `__regex_action_free:n`, `__regex_action_free_group:n`, and `__regex_action_free_aux:nn`.)

`__regex_action_cost:n` A transition which consumes the current character and shifts the state by #1. The resulting state is stored in the appropriate array for use at the next position, and we also store the current submatches.

```

26282 \cs_new_protected:Npn \__regex_action_cost:n #1
26283 {
26284   \exp_args:Nx \__regex_store_state:n
26285   { \int_eval:n { \l__regex_curr_state_int + #1 } }
26286 }

```

(End definition for __regex_action_cost:n.)

__regex_store_state:n Put the given state and current submatch information in \g__regex_thread_info_intarray, and increment the length of the array.

```

26287 \cs_new_protected:Npn \__regex_store_state:n #1
26288 {
26289   \exp_args:No \__regex_store_submatches:nn
26290   \l__regex_curr_submatches_tl {#1}
26291   \int_incr:N \l__regex_max_thread_int
26292 }
26293 \cs_new_protected:Npn \__regex_store_submatches:nn #1#2
26294 {
26295   \__kernel_intarray_gset_range_from_clist:Nnn
26296   \g__regex_thread_info_intarray
26297   {
26298     \__regex_int_eval:w
26299     1 + \l__regex_max_thread_int *
26300     (\l__regex_capturing_group_int * 2 + 1)
26301   }
26302   { #2 , #1 }
26303 }

```

(End definition for __regex_store_state:n and __regex_store_submatches:.)

__regex_disable_submatches: Some user functions don't require tracking submatches. We get a performance improvement by simply defining the relevant functions to remove their argument and do nothing with it.

```

26304 \cs_new_protected:Npn \__regex_disable_submatches:
26305 {
26306   \cs_set_protected:Npn \__regex_store_submatches:n ##1 { }
26307   \cs_set_protected:Npn \__regex_action_submatch:nN ##1##2 { }
26308 }

```

(End definition for __regex_disable_submatches:.)

__regex_action_submatch:nN Update the current submatches with the information from the current position. Maybe a bottleneck.

```

\__regex_action_submatch_aux:w
\__regex_action_submatch_auxii:w
\__regex_action_submatch_auxiii:w
\__regex_action_submatch_auxiv:w
26309 \cs_new_protected:Npn \__regex_action_submatch:nN #1#2
26310 {
26311   \exp_after:wN \__regex_action_submatch_aux:w
26312   \l__regex_curr_submatches_tl ; {#1} #2
26313 }
26314 \cs_new_protected:Npn \__regex_action_submatch_aux:w #1 ; #2#3
26315 {
26316   \tl_set:Nx \l__regex_curr_submatches_tl
26317   {
26318     \prg_replicate:nn
26319     { #2 \if_meaning:w > #3 + \l__regex_capturing_group_int \fi: }

```

```

26320         { \__regex_action_submatch_auxii:w }
26321         \__regex_action_submatch_auxiii:w
26322         #1
26323     }
26324 }
26325 \cs_new:Npn \__regex_action_submatch_auxii:w
26326     #1 \__regex_action_submatch_auxiii:w #2 ,
26327     { #2 , #1 \__regex_action_submatch_auxiii:w }
26328 \cs_new:Npn \__regex_action_submatch_auxiii:w #1 ,
26329     { \int_use:N \l__regex_curr_pos_int , }

```

(End definition for __regex_action_submatch:nN and others.)

__regex_action_success: There is a successful match when an execution path reaches the last state in the NFA, unless this marks a second identical empty match. Then mark that there was a successful match; it is empty if it is “fresh”; and we store the current position and submatches. The current step is then interrupted with \prg_break:, and only paths with higher precedence are pursued further. The values stored here may be overwritten by a later success of a path with higher precedence.

```

26330 \cs_new_protected:Npn \__regex_action_success:
26331 {
26332     \__regex_if_two_empty_matches:F
26333     {
26334         \bool_set_true:N \l__regex_match_success_bool
26335         \bool_set_eq:NN \l__regex_empty_success_bool
26336         \l__regex_fresh_thread_bool
26337         \int_set_eq:NN \l__regex_success_pos_int \l__regex_curr_pos_int
26338         \int_set_eq:NN \l__regex_last_char_success_int \l__regex_last_char_int
26339         \tl_build_clear:N \l__regex_matched_analysis_tl
26340         \tl_set_eq:NN \l__regex_success_submatches_tl
26341         \l__regex_curr_submatches_tl
26342         \prg_break:
26343     }
26344 }

```

(End definition for __regex_action_success:.)

41.6 Replacement

41.6.1 Variables and helpers used in replacement

\l__regex_replacement_csnames_int The behaviour of closing braces inside a replacement text depends on whether a sequences \c{ or \u{ has been encountered. The number of “open” such sequences that should be closed by } is stored in \l__regex_replacement_csnames_int, and decreased by 1 by each }.

```

26345 \int_new:N \l__regex_replacement_csnames_int

```

(End definition for \l__regex_replacement_csnames_int.)

\l__regex_replacement_category_tl This sequence of letters is used to correctly restore categories in nested constructions such as \cL(abc\cD(_)d).

\l__regex_replacement_category_seq

```

26346 \tl_new:N \l__regex_replacement_category_tl
26347 \seq_new:N \l__regex_replacement_category_seq

```

(End definition for \l__regex_replacement_category_tl and \l__regex_replacement_category_seq.)

\l__regex_balance_tl This token list holds the replacement text for __regex_replacement_balance_one_match:n while it is being built incrementally.

26348 \tl_new:N \l__regex_balance_tl

(End definition for \l__regex_balance_tl.)

__regex_replacement_balance_one_match:n This expects as an argument the first index of a set of entries in \g__regex_submatch_begin_intarray (and related arrays) which hold the submatch information for a given match. It can be used within an integer expression to obtain the brace balance incurred by performing the replacement on that match. This combines the braces lost by removing the match, braces added by all the submatches appearing in the replacement, and braces appearing explicitly in the replacement. Even though it is always redefined before use, we initialize it as for an empty replacement. An important property is that concatenating several calls to that function must result in a valid integer expression (hence a leading + in the actual definition).

26349 \cs_new:Npn __regex_replacement_balance_one_match:n #1

26350 { - __regex_submatch_balance:n {#1} }

(End definition for __regex_replacement_balance_one_match:n.)

__regex_replacement_do_one_match:n The input is the same as __regex_replacement_balance_one_match:n. This function is redefined to expand to the part of the token list from the end of the previous match to a given match, followed by the replacement text. Hence concatenating the result of this function with all possible arguments (one call for each match), as well as the range from the end of the last match to the end of the string, produces the fully replaced token list. The initialization does not matter, but (as an example) we set it as for an empty replacement.

26351 \cs_new:Npn __regex_replacement_do_one_match:n #1

26352 {

26353 __regex_query_range:nn

26354 { __kernel_intarray_item:Nn \g__regex_submatch_prev_intarray {#1} }

26355 { __kernel_intarray_item:Nn \g__regex_submatch_begin_intarray {#1} }

26356 }

(End definition for __regex_replacement_do_one_match:n.)

__regex_replacement_exp_not:N This function lets us navigate around the fact that the primitive \exp_not:n requires a braced argument. As far as I can tell, it is only needed if the user tries to include in the replacement text a control sequence set equal to a macro parameter character, such as \c_parameter_token. Indeed, within an x-expanding assignment, \exp_not:N # behaves as a single #, whereas \exp_not:n {#} behaves as a doubled ##.

26357 \cs_new:Npn __regex_replacement_exp_not:N #1 { \exp_not:n {#1} }

(End definition for __regex_replacement_exp_not:N.)

__regex_replacement_exp_not:V This is used for the implementation of \u, and it gets redefined for \peek_regex_replace_once:nnTF.

26358 \cs_new_eq:NN __regex_replacement_exp_not:V \exp_not:V

(End definition for __regex_replacement_exp_not:V.)

41.6.2 Query and brace balance

`__regex_query_range:nn` When it is time to extract submatches from the token list, the various tokens are stored in `\toks` registers numbered from `\l__regex_min_pos_int` inclusive to `\l__regex_max_pos_int` exclusive. The function `__regex_query_range:nn {<min>} {<max>}` unpacks registers from the position `<min>` to the position `<max> - 1` included. Once this is expanded, a second x-expansion results in the actual tokens from the query. That second expansion is only done by user functions at the very end of their operation, after checking (and correcting) the brace balance first.

```

26359 \cs_new:Npn \__regex_query_range:nn #1#2
26360 {
26361   \exp_after:wN \__regex_query_range_loop:ww
26362   \int_value:w \__regex_int_eval:w #1 \exp_after:wN ;
26363   \int_value:w \__regex_int_eval:w #2 ;
26364   \prg_break_point:
26365 }
26366 \cs_new:Npn \__regex_query_range_loop:ww #1 ; #2 ;
26367 {
26368   \if_int_compare:w #1 < #2 \exp_stop_f:
26369   \else:
26370     \exp_after:wN \prg_break:
26371   \fi:
26372   \__regex_toks_use:w #1 \exp_stop_f:
26373   \exp_after:wN \__regex_query_range_loop:ww
26374   \int_value:w \__regex_int_eval:w #1 + 1 ; #2 ;
26375 }
```

(End definition for `__regex_query_range:nn` and `__regex_query_range_loop:ww`.)

`__regex_query_submatch:n` Find the start and end positions for a given submatch (of a given match).

```

26376 \cs_new:Npn \__regex_query_submatch:n #1
26377 {
26378   \__regex_query_range:nn
26379   { \__kernel_intarray_item:Nn \g__regex_submatch_begin_intarray {#1} }
26380   { \__kernel_intarray_item:Nn \g__regex_submatch_end_intarray {#1} }
26381 }
```

(End definition for `__regex_query_submatch:n`.)

`__regex_submatch_balance:n` Every user function must result in a balanced token list (unbalanced token lists cannot be stored by TeX). When we unpacked the query, we kept track of the brace balance, hence the contribution from a given range is the difference between the brace balances at the `<max pos>` and `<min pos>`. These two positions are found in the corresponding “submatch” arrays.

```

26382 \cs_new_protected:Npn \__regex_submatch_balance:n #1
26383 {
26384   \int_eval:n
26385   {
26386     \__regex_intarray_item:NnF \g__regex_balance_intarray
26387     {
26388       \__kernel_intarray_item:Nn
26389       \g__regex_submatch_end_intarray {#1}
26390     }
26391     { 0 }
```

```

26392 -
26393 \__regex_intarray_item:NnF \g__regex_balance_intarray
26394 {
26395     \__kernel_intarray_item:Nn
26396     \g__regex_submatch_begin_intarray {#1}
26397 }
26398 { 0 }
26399 }
26400 }

```

(End definition for __regex_submatch_balance:n.)

41.6.3 Framework

```

\__regex_replacement:n
\__regex_replacement_aux:n

```

The replacement text is built incrementally. We keep track in \l__regex_balance_int of the balance of explicit begin- and end-group tokens and we store in \l__regex_balance_tl some code to compute the brace balance from submatches (see its description). Detect unescaped right braces, and escaped characters, with trailing \prg_do_nothing: because some of the later function look-ahead. Once the whole replacement text has been parsed, make sure that there is no open csname. Finally, define the balance_one_match and do_one_match functions.

```

26401 \cs_new_protected:Npn \__regex_replacement:n #1
26402 {
26403     \group_begin:
26404     \tl_build_begin:N \l__regex_build_tl
26405     \int_zero:N \l__regex_balance_int
26406     \tl_clear:N \l__regex_balance_tl
26407     \__regex_escape_use:nnnn
26408     {
26409         \if_charcode:w \c_right_brace_str ##1
26410             \__regex_replacement_rbrace:N
26411         \else:
26412             \__regex_replacement_normal:n
26413         \fi:
26414         ##1
26415     }
26416     { \__regex_replacement_escaped:N ##1 }
26417     { \__regex_replacement_normal:n ##1 }
26418     {#1}
26419     \prg_do_nothing: \prg_do_nothing:
26420     \if_int_compare:w \l__regex_replacement_csnames_int > 0 \exp_stop_f:
26421         \__kernel_msg_error:nnx { kernel } { replacement-missing-rbrace }
26422         { \int_use:N \l__regex_replacement_csnames_int }
26423         \tl_build_put_right:Nx \l__regex_build_tl
26424         { \prg_replicate:nn \l__regex_replacement_csnames_int \cs_end: }
26425     \fi:
26426     \seq_if_empty:NF \l__regex_replacement_category_seq
26427     {
26428         \__kernel_msg_error:nnx { kernel } { replacement-missing-rparen }
26429         { \seq_count:N \l__regex_replacement_category_seq }
26430         \seq_clear:N \l__regex_replacement_category_seq
26431     }
26432     \cs_gset:Npx \__regex_replacement_balance_one_match:n ##1
26433     {

```

```

26434         + \int_use:N \l__regex_balance_int
26435         \l__regex_balance_tl
26436         - \__regex_submatch_balance:n {##1}
26437     }
26438     \tl_build_end:N \l__regex_build_tl
26439     \exp_args:NNo
26440     \group_end:
26441     \__regex_replacement_aux:n \l__regex_build_tl
26442 }
26443 \cs_new_protected:Npn \__regex_replacement_aux:n #1
26444 {
26445     \cs_set:Npn \__regex_replacement_do_one_match:n ##1
26446     {
26447         \__regex_query_range:nn
26448         {
26449             \__kernel_intarray_item:Nn
26450             \g__regex_submatch_prev_intarray {##1}
26451         }
26452         {
26453             \__kernel_intarray_item:Nn
26454             \g__regex_submatch_begin_intarray {##1}
26455         }
26456         #1
26457     }
26458 }

```

(End definition for __regex_replacement:n and __regex_replacement_aux:n.)

__regex_replacement_put:n This gets redefined for \peek_regex_replace_once:nnTF.

```

26459 \cs_new_protected:Npn \__regex_replacement_put:n
26460 { \tl_build_put_right:Nn \l__regex_build_tl }

```

(End definition for __regex_replacement_put:n.)

__regex_replacement_normal:n Most characters are simply sent to the output by \tl_build_put_right:Nn, unless a particular category code has been requested: then __regex_replacement_c_A:w or a similar auxiliary is called. One exception is right parentheses, which restore the category code in place before the group started. Note that the sequence is non-empty there: it contains an empty entry corresponding to the initial value of \l__regex_replacement_category_tl. The argument #1 can be a space, otherwise it is a single character.

```

26461 \cs_new_protected:Npn \__regex_replacement_normal:n #1
26462 {
26463     \tl_if_empty:NTF \l__regex_replacement_category_tl
26464     { \__regex_replacement_put:n {#1} }
26465     { % (
26466         \token_if_eq_charcode:NNTF #1 )
26467         {
26468             \seq_pop:NN \l__regex_replacement_category_seq
26469             \l__regex_replacement_category_tl
26470         }
26471         {
26472             \use:c
26473             {
26474                 __regex_replacement_c_

```

```

26475         \l__regex_replacement_category_tl :w
26476     }
26477     \__regex_replacement_normal:n {#1}
26478 }
26479 }
26480 }

```

(End definition for __regex_replacement_normal:n.)

__regex_replacement_escaped:N As in parsing a regular expression, we use an auxiliary built from #1 if defined. Otherwise, check for escaped digits (standing from submatches from 0 to 9): anything else is a raw character. We use \token_to_str:N to give spaces the right category code.

```

26481 \cs_new_protected:Npn \__regex_replacement_escaped:N #1
26482 {
26483     \cs_if_exist_use:cF { __regex_replacement_#1:w }
26484     {
26485         \if_int_compare:w 1 < 1#1 \exp_stop_f:
26486         \__regex_replacement_put_submatch:n {#1}
26487     \else:
26488         \exp_args:No \__regex_replacement_normal:n
26489         { \token_to_str:N #1 }
26490     \fi:
26491 }
26492 }

```

(End definition for __regex_replacement_escaped:N.)

41.6.4 Submatches

__regex_replacement_put_submatch:n Insert a submatch in the replacement text. This is dropped if the submatch number is larger than the number of capturing groups. Unless the submatch appears inside a \c{...} or \u{...} construction, it must be taken into account in the brace balance. Later on, ##1 will be replaced by a pointer to the 0-th submatch for a given match. There is an \exp_not:N here as at the point-of-use of \l__regex_balance_tl there is an x-type expansion which is needed to get ##1 in correctly.

```

26493 \cs_new_protected:Npn \__regex_replacement_put_submatch:n #1
26494 {
26495     \if_int_compare:w #1 < \l__regex_capturing_group_int
26496     \__regex_replacement_put_submatch_aux:n {#1}
26497     \fi:
26498 }
26499 \cs_new_protected:Npn \__regex_replacement_put_submatch_aux:n #1
26500 {
26501     \tl_build_put_right:Nn \l__regex_build_tl
26502     { \__regex_query_submatch:n { \int_eval:n { #1 + ##1 } } }
26503     \if_int_compare:w \l__regex_replacement_csnames_int = 0 \exp_stop_f:
26504     \tl_put_right:Nn \l__regex_balance_tl
26505     {
26506         + \__regex_submatch_balance:n
26507         { \exp_not:N \int_eval:n { #1 + ##1 } }
26508     }
26509     \fi:
26510 }

```

(End definition for `__regex_replacement_put_submatch:n` and `__regex_replacement_put_submatch_aux:n`.)

`__regex_replacement_g:w` Grab digits for the `\g` escape sequence in a primitive assignment to the integer `\l__regex_internal_a_int`. At the end of the run of digits, check that it ends with a right brace.

```

26511 \cs_new_protected:Npn __regex_replacement_g:w #1#2
26512 {
26513   __regex_two_if_eq:NNNTF
26514   #1 #2 __regex_replacement_normal:n \c_left_brace_str
26515   { \l__regex_internal_a_int = __regex_replacement_g_digits:NN }
26516   { __regex_replacement_error:NNN g #1 #2 }
26517 }
26518 \cs_new:Npn __regex_replacement_g_digits:NN #1#2
26519 {
26520   \token_if_eq_meaning:NNTF #1 __regex_replacement_normal:n
26521   {
26522     \if_int_compare:w 1 < 1#2 \exp_stop_f:
26523     #2
26524     \exp_after:wN \use_i:nnn
26525     \exp_after:wN __regex_replacement_g_digits:NN
26526   \else:
26527     \exp_stop_f:
26528     \exp_after:wN __regex_replacement_error:NNN
26529     \exp_after:wN g
26530   \fi:
26531 }
26532 {
26533   \exp_stop_f:
26534   \if_meaning:w __regex_replacement_rbrace:N #1
26535   \exp_args:No __regex_replacement_put_submatch:n
26536   { \int_use:N \l__regex_internal_a_int }
26537   \exp_after:wN \use_none:nn
26538   \else:
26539     \exp_after:wN __regex_replacement_error:NNN
26540     \exp_after:wN g
26541   \fi:
26542 }
26543 #1 #2
26544 }
```

(End definition for `__regex_replacement_g:w` and `__regex_replacement_g_digits:NN`.)

41.6.5 Csnames in replacement

`__regex_replacement_c:w` `\c` may only be followed by an unescaped character. If followed by a left brace, start a control sequence by calling an auxiliary common with `\u`. Otherwise test whether the category is known; if it is not, complain.

```

26545 \cs_new_protected:Npn __regex_replacement_c:w #1#2
26546 {
26547   \token_if_eq_meaning:NNTF #1 __regex_replacement_normal:n
26548   {
26549     \exp_after:wN \token_if_eq_charcode:NNTF \c_left_brace_str #2
26550     { __regex_replacement_cu_aux:Nw __regex_replacement_exp_not:N }
```

```

26551         {
26552             \cs_if_exist:cTF { __regex_replacement_c_#2:w }
26553             { \__regex_replacement_cat:NNN #2 }
26554             { \__regex_replacement_error:NNN c #1#2 }
26555         }
26556     }
26557     { \__regex_replacement_error:NNN c #1#2 }
26558 }

```

(End definition for __regex_replacement_c:w.)

__regex_replacement_cu_aux:Nw Start a control sequence with \cs:w, protected from expansion by #1 (either __regex_replacement_exp_not:N or \exp_not:V), or turned to a string by \tl_to_str:V if inside another csname construction \c or \u. We use \tl_to_str:V rather than \tl_to_str:N to deal with integers and other registers.

```

26559 \cs_new_protected:Npn \__regex_replacement_cu_aux:Nw #1
26560 {
26561     \if_case:w \l__regex_replacement_csnames_int
26562     \tl_build_put_right:Nn \l__regex_build_tl
26563     { \exp_not:n { \exp_after:wN #1 \cs:w } }
26564     \else:
26565     \tl_build_put_right:Nn \l__regex_build_tl
26566     { \exp_not:n { \exp_after:wN \tl_to_str:V \cs:w } }
26567     \fi:
26568     \int_incr:N \l__regex_replacement_csnames_int
26569 }

```

(End definition for __regex_replacement_cu_aux:Nw.)

__regex_replacement_u:w Check that \u is followed by a left brace. If so, start a control sequence with \cs:w, which is then unpacked either with \exp_not:V or \tl_to_str:V depending on the current context.

```

26570 \cs_new_protected:Npn \__regex_replacement_u:w #1#2
26571 {
26572     \__regex_two_if_eq:NNNTF
26573     #1 #2 \__regex_replacement_normal:n \c_left_brace_str
26574     { \__regex_replacement_cu_aux:Nw \__regex_replacement_exp_not:V }
26575     { \__regex_replacement_error:NNN u #1#2 }
26576 }

```

(End definition for __regex_replacement_u:w.)

__regex_replacement_rbrace:N Within a \c{...} or \u{...} construction, end the control sequence, and decrease the brace count. Otherwise, this is a raw right brace.

```

26577 \cs_new_protected:Npn \__regex_replacement_rbrace:N #1
26578 {
26579     \if_int_compare:w \l__regex_replacement_csnames_int > 0 \exp_stop_f:
26580     \tl_build_put_right:Nn \l__regex_build_tl { \cs_end: }
26581     \int_decr:N \l__regex_replacement_csnames_int
26582     \else:
26583     \__regex_replacement_normal:n {#1}
26584     \fi:
26585 }

```

(End definition for __regex_replacement_rbrace:N.)

41.6.6 Characters in replacement

`__regex_replacement_cat:NNN` Here, `#1` is a letter among BEMTPUDSLOA and `#2#3` denote the next character. Complain if we reach the end of the replacement or if the construction appears inside `\c{...}` or `\u{...}`, and detect the case of a parenthesis. In that case, store the current category in a sequence and switch to a new one.

```

26586 \cs_new_protected:Npn __regex_replacement_cat:NNN #1#2#3
26587 {
26588   \token_if_eq_meaning:NNTF \prg_do_nothing: #3
26589   { __kernel_msg_error:nn { kernel } { replacement-catcode-end } }
26590   {
26591     \int_compare:nNnTF { \l__regex_replacement_csnames_int } > 0
26592     {
26593       __kernel_msg_error:nnnn
26594       { kernel } { replacement-catcode-in-cs } {#1} {#3}
26595       #2 #3
26596     }
26597     {
26598       __regex_two_if_eq:NNNNTF #2 #3 __regex_replacement_normal:n (
26599       {
26600         \seq_push:NV \l__regex_replacement_category_seq
26601         \l__regex_replacement_category_tl
26602         \tl_set:Nn \l__regex_replacement_category_tl {#1}
26603       }
26604       {
26605         \token_if_eq_meaning:NNT #2 __regex_replacement_escaped:N
26606         {
26607           __regex_char_if_alphanumeric:NTF #3
26608           {
26609             __kernel_msg_error:nnnn
26610             { kernel } { replacement-catcode-escaped }
26611             {#1} {#3}
26612           }
26613           { }
26614         }
26615         \use:c { __regex_replacement_c_#1:w } #2 #3
26616       }
26617     }
26618   }
26619 }
```

(End definition for `__regex_replacement_cat:NNN`.)

We now need to change the category code of the null character many times, hence work in a group. The catcode-specific macros below are defined in alphabetical order; if you are trying to understand the code, start from the end of the alphabet as those categories are simpler than active or begin-group.

```

26620 \group_begin:
```

`__regex_replacement_char:nNN` The only way to produce an arbitrary character-catcode pair is to use the `\lowercase` or `\uppercase` primitives. This is a wrapper for our purposes. The first argument is the null character with various catcodes. The second and third arguments are grabbed from the input stream: `#3` is the character whose character code to reproduce. We could use

`\char_generate:nn` but only for some catcodes (active characters and spaces are not supported).

```

26621 \cs_new_protected:Npn \__regex_replacement_char:nNN #1#2#3
26622 {
26623   \tex_lccode:D 0 = '#3 \scan_stop:
26624   \tex_lowercase:D { \__regex_replacement_put:n {#1} }
26625 }

```

(End definition for `__regex_replacement_char:nNN`.)

`__regex_replacement_c_A:w` For an active character, expansion must be avoided, twice because we later do two x-expansions, to unpack `\toks` for the query, and to expand their contents to tokens of the query.

```

26626 \char_set_catcode_active:N \^^@
26627 \cs_new_protected:Npn \__regex_replacement_c_A:w
26628 { \__regex_replacement_char:nNN { \exp_not:n { \exp_not:N \^^@ } } }

```

(End definition for `__regex_replacement_c_A:w`.)

`__regex_replacement_c_B:w` An explicit begin-group token increases the balance, unless within a `\c{...}` or `\u{...}` construction. Add the desired begin-group character, using the standard `\if_false:` trick. We eventually x-expand twice. The first time must yield a balanced token list, and the second one gives the bare begin-group token. The `\exp_after:wN` is not strictly needed, but is more consistent with l3tl-analysis.

```

26629 \char_set_catcode_group_begin:N \^^@
26630 \cs_new_protected:Npn \__regex_replacement_c_B:w
26631 {
26632   \if_int_compare:w \l__regex_replacement_csnames_int = 0 \exp_stop_f:
26633   \int_incr:N \l__regex_balance_int
26634   \fi:
26635   \__regex_replacement_char:nNN
26636   { \exp_not:n { \exp_after:wN \^^@ \if_false: } \fi: } }
26637 }

```

(End definition for `__regex_replacement_c_B:w`.)

`__regex_replacement_c_C:w` This is not quite catcode-related: when the user requests a character with category “control sequence”, the one-character control symbol is returned. As for the active character, we prepare for two x-expansions.

```

26638 \cs_new_protected:Npn \__regex_replacement_c_C:w #1#2
26639 {
26640   \tl_build_put_right:Nn \l__regex_build_tl
26641   { \exp_not:N \__regex_replacement_exp_not:N \exp_not:c {#2} }
26642 }

```

(End definition for `__regex_replacement_c_C:w`.)

`__regex_replacement_c_D:w` Subscripts fit the mould: `\lowercase` the null byte with the correct category.

```

26643 \char_set_catcode_math_subscript:N \^^@
26644 \cs_new_protected:Npn \__regex_replacement_c_D:w
26645 { \__regex_replacement_char:nNN { \^^@ } }

```

(End definition for `__regex_replacement_c_D:w`.)

`_regex_replacement_c_E:w` Similar to the begin-group case, the second x-expansion produces the bare end-group token.

```

26646 \char_set_catcode_group_end:N \^^@
26647 \cs_new_protected:Npn \_regex_replacement_c_E:w
26648 {
26649   \if_int_compare:w \l__regex_replacement_csnames_int = 0 \exp_stop_f:
26650   \int_decr:N \l__regex_balance_int
26651   \fi:
26652   \_regex_replacement_char:nNN
26653   { \exp_not:n { \if_false: { \fi: ^^@ } }
26654   }

```

(End definition for _regex_replacement_c_E:w.)

`_regex_replacement_c_L:w` Simply `\lowercase` a letter null byte to produce an arbitrary letter.

```

26655 \char_set_catcode_letter:N \^^@
26656 \cs_new_protected:Npn \_regex_replacement_c_L:w
26657 { \_regex_replacement_char:nNN { ^^@ } }

```

(End definition for _regex_replacement_c_L:w.)

`_regex_replacement_c_M:w` No surprise here, we lowercase the null math toggle.

```

26658 \char_set_catcode_math_toggle:N \^^@
26659 \cs_new_protected:Npn \_regex_replacement_c_M:w
26660 { \_regex_replacement_char:nNN { ^^@ } }

```

(End definition for _regex_replacement_c_M:w.)

`_regex_replacement_c_O:w` Lowercase an other null byte.

```

26661 \char_set_catcode_other:N \^^@
26662 \cs_new_protected:Npn \_regex_replacement_c_O:w
26663 { \_regex_replacement_char:nNN { ^^@ } }

```

(End definition for _regex_replacement_c_O:w.)

`_regex_replacement_c_P:w` For macro parameters, expansion is a tricky issue. We need to prepare for two x-expansions and passing through various macro definitions. Note that we cannot replace one `\exp_not:n` by doubling the macro parameter characters because this would misbehave if a mischievous user asks for `\c{\cP\#}`, since that macro parameter character would be doubled.

```

26664 \char_set_catcode_parameter:N \^^@
26665 \cs_new_protected:Npn \_regex_replacement_c_P:w
26666 {
26667   \_regex_replacement_char:nNN
26668   { \exp_not:n { \exp_not:n { ^^@^^@^^@^^@ } } }
26669 }

```

(End definition for _regex_replacement_c_P:w.)

`_regex_replacement_c_S:w` Spaces are normalized on input by \TeX to have character code 32. It is in fact impossible to get a token with character code 0 and category code 10. Hence we use 32 instead of 0 as our base character.

```

26670 \cs_new_protected:Npn \_regex_replacement_c_S:w #1#2
26671 {

```

```

26672 \if_int_compare:w '#2 = 0 \exp_stop_f:
26673 \__kernel_msg_error:nn { kernel } { replacement-null-space }
26674 \fi:
26675 \tex_lccode:D '\ = '#2 \scan_stop:
26676 \tex_lowercase:D { \__regex_replacement_put:n {~} }
26677 }

```

(End definition for __regex_replacement_c_S:w.)

__regex_replacement_c_T:w No surprise for alignment tabs here. Those are surrounded by the appropriate braces whenever necessary, hence they don't cause trouble in alignment settings.

```

26678 \char_set_catcode_alignment:N \^^@
26679 \cs_new_protected:Npn \__regex_replacement_c_T:w
26680 { \__regex_replacement_char:nnn { ^^@ } }

```

(End definition for __regex_replacement_c_T:w.)

__regex_replacement_c_U:w Simple call to __regex_replacement_char:nnn which lowercases the math superscript ^^@.

```

26681 \char_set_catcode_math_superscript:N \^^@
26682 \cs_new_protected:Npn \__regex_replacement_c_U:w
26683 { \__regex_replacement_char:nnn { ^^@ } }

```

(End definition for __regex_replacement_c_U:w.)

Restore the catcode of the null byte.

```

26684 \group_end:

```

41.6.7 An error

_regex_replacement_error:NNN Simple error reporting by calling one of the messages replacement-c, replacement-g, or replacement-u.

```

26685 \cs_new_protected:Npn \_regex_replacement_error:NNN #1#2#3
26686 {
26687 \__kernel_msg_error:nnx { kernel } { replacement-#1 } {#3}
26688 #2 #3
26689 }

```

(End definition for _regex_replacement_error:NNN.)

41.7 User functions

\regex_new:N Before being assigned a sensible value, a regex variable matches nothing.

```

26690 \cs_new_protected:Npn \regex_new:N #1
26691 { \cs_new_eq:NN #1 \c__regex_no_match_regex }

```

(End definition for \regex_new:N. This function is documented on page 236.)

\l_tmpa_regex The usual scratch space.

```

\l_tmpb_regex 26692 \regex_new:N \l_tmpa_regex
\g_tmpa_regex 26693 \regex_new:N \l_tmpb_regex
\g_tmpb_regex 26694 \regex_new:N \g_tmpa_regex
26695 \regex_new:N \g_tmpb_regex

```

(End definition for \l_tmpa_regex and others. These variables are documented on page 238.)

\regex_set:Nn Compile, then store the result in the user variable with the appropriate assignment function.
\regex_gset:Nn
\regex_const:Nn

```

26696 \cs_new_protected:Npn \regex_set:Nn #1#2
26697 {
26698   \__regex_compile:n {#2}
26699   \tl_set_eq:NN #1 \l__regex_internal_regex
26700 }
26701 \cs_new_protected:Npn \regex_gset:Nn #1#2
26702 {
26703   \__regex_compile:n {#2}
26704   \tl_gset_eq:NN #1 \l__regex_internal_regex
26705 }
26706 \cs_new_protected:Npn \regex_const:Nn #1#2
26707 {
26708   \__regex_compile:n {#2}
26709   \tl_const:Nx #1 { \exp_not:o \l__regex_internal_regex }
26710 }

```

(End definition for \regex_set:Nn, \regex_gset:Nn, and \regex_const:Nn. These functions are documented on page 236.)

\regex_show:N User functions: the n variant requires compilation first. Then show the variable with
\regex_show:n some appropriate text. The auxiliary is defined in a different section.

```

26711 \cs_new_protected:Npn \regex_show:n #1
26712 {
26713   \__regex_compile:n {#1}
26714   \__regex_show:N \l__regex_internal_regex
26715   \msg_show:nnxxxx { LaTeX / kernel } { show-regex }
26716   { \tl_to_str:n {#1} } { }
26717   { \l__regex_internal_a_tl } { }
26718 }
26719 \cs_new_protected:Npn \regex_show:N #1
26720 {
26721   \__kernel_chk_defined:NT #1
26722   {
26723     \__regex_show:N #1
26724     \msg_show:nnxxxx { LaTeX / kernel } { show-regex }
26725     { } { \token_to_str:N #1 }
26726     { \l__regex_internal_a_tl } { }
26727   }
26728 }

```

(End definition for \regex_show:N and \regex_show:n. These functions are documented on page 236.)

\regex_match:nnTF Those conditionals are based on a common auxiliary defined later. Its first argument
\regex_match:NnTF builds the NFA corresponding to the regex, and the second argument is the query token list. Once we have performed the match, convert the resulting boolean to \prg_return_true: or false.

```

26729 \prg_new_protected_conditional:Npnn \regex_match:nn #1#2 { T , F , TF }
26730 {
26731   \__regex_if_match:nn { \__regex_build:n {#1} } {#2}
26732   \__regex_return:
26733 }
26734 \prg_new_protected_conditional:Npnn \regex_match:Nn #1#2 { T , F , TF }

```

```

26735 {
26736   \__regex_if_match:nn { \__regex_build:N #1 } {#2}
26737   \__regex_return:
26738 }

```

(End definition for \regex_match:nnTF and \regex_match:NnTF. These functions are documented on page 236.)

\regex_count:nnN Again, use an auxiliary whose first argument builds the NFA.

```

\regex_count:NnN
26739 \cs_new_protected:Npn \regex_count:nnN #1
26740 { \__regex_count:nnN { \__regex_build:n {#1} } }
26741 \cs_new_protected:Npn \regex_count:NnN #1
26742 { \__regex_count:nnN { \__regex_build:N #1 } }

```

(End definition for \regex_count:nnN and \regex_count:NnN. These functions are documented on page 237.)

\regex_extract_once:nnN We define here 40 user functions, following a common pattern in terms of :nnN auxiliaries, defined in the coming subsections. The auxiliary is handed __regex_build:n or __regex_build:N with the appropriate regex argument, then all other necessary arguments (replacement text, token list, etc. The conditionals call __regex_return: to return either true or false once matching has been performed.

```

\regex_extract_once:nnN
\regex_extract_once:nnNTF
\regex_extract_once:NnN
\regex_extract_once:NnNTF
\regex_extract_all:nnN
\regex_extract_all:nnNTF
\regex_extract_all:NnN
\regex_extract_all:NnNTF
\regex_replace_once:nnN
\regex_replace_once:nnNTF
\regex_replace_once:NnN
\regex_replace_once:NnNTF
\regex_replace_all:nnN
\regex_replace_all:nnNTF
\regex_replace_all:NnN
\regex_replace_all:NnNTF
\regex_split:nnN
\regex_split:nnNTF
\regex_split:NnN
\regex_split:NnNTF
26743 \cs_set_protected:Npn \__regex_tmp:w #1#2#3
26744 {
26745   \cs_new_protected:Npn #2 ##1 { #1 { \__regex_build:n {##1} } }
26746   \cs_new_protected:Npn #3 ##1 { #1 { \__regex_build:N ##1 } }
26747   \prg_new_protected_conditional:Npnn #2 ##1##2##3 { T , F , TF }
26748     { #1 { \__regex_build:n {##1} } {##2} ##3 \__regex_return: }
26749   \prg_new_protected_conditional:Npnn #3 ##1##2##3 { T , F , TF }
26750     { #1 { \__regex_build:N ##1 } {##2} ##3 \__regex_return: }
26751 }
26752 \__regex_tmp:w \__regex_extract_once:nnN
26753 \__regex_tmp:w \__regex_extract_once:nnNTF
26754 \__regex_tmp:w \__regex_extract_all:nnN
26755 \__regex_tmp:w \__regex_extract_all:nnNTF
26756 \__regex_tmp:w \__regex_replace_once:nnN
26757 \__regex_tmp:w \__regex_replace_once:nnNTF
26758 \__regex_tmp:w \__regex_replace_all:nnN
26759 \__regex_tmp:w \__regex_replace_all:nnNTF
26760 \__regex_tmp:w \__regex_split:nnN
\__regex_split:nnNTF
\__regex_split:NnN
\__regex_split:NnNTF

```

(End definition for \regex_extract_once:nnN and others. These functions are documented on page 237.)

41.7.1 Variables and helpers for user functions

\l__regex_match_count_int The number of matches found so far is stored in \l__regex_match_count_int. This is only used in the \regex_count:nnN functions.

```

26761 \int_new:N \l__regex_match_count_int

```

(End definition for \l__regex_match_count_int.)

__regex_begin Those flags are raised to indicate extra begin-group or end-group tokens when extracting submatches.

```

26762 \flag_new:n { \__regex_begin }
26763 \flag_new:n { \__regex_end }

```

(End definition for `__regex_begin` and `__regex_end`.)

`\l__regex_min_submatch_int` The end-points of each submatch are stored in two arrays whose index *<submatch>* ranges from `\l__regex_min_submatch_int` (inclusive) to `\l__regex_submatch_int` (exclusive). Each successful match comes with a 0-th submatch (the full match), and one match for each capturing group: submatches corresponding to the last successful match are labelled starting at `zeroth_submatch`. The entry `\l__regex_zeroth_submatch_int` in `\g__regex_submatch_prev_intarray` holds the position at which that match attempt started: this is used for splitting and replacements.

```
26764 \int_new:N \l__regex_min_submatch_int
26765 \int_new:N \l__regex_submatch_int
26766 \int_new:N \l__regex_zeroth_submatch_int
```

(End definition for `\l__regex_min_submatch_int`, `\l__regex_submatch_int`, and `\l__regex_zeroth_submatch_int`.)

`\g__regex_submatch_prev_intarray` Hold the place where the match attempt begun and the end-points of each submatch.

```
\g__regex_submatch_begin_intarray 26767 \intarray_new:Nn \g__regex_submatch_prev_intarray { 65536 }
\g__regex_submatch_end_intarray    26768 \intarray_new:Nn \g__regex_submatch_begin_intarray { 65536 }
                                   26769 \intarray_new:Nn \g__regex_submatch_end_intarray { 65536 }
```

(End definition for `\g__regex_submatch_prev_intarray`, `\g__regex_submatch_begin_intarray`, and `\g__regex_submatch_end_intarray`.)

`\g__regex_balance_intarray` The first thing we do when matching is to store the balance of begin-group/end-group characters into `\g__regex_balance_intarray`.

```
26770 \intarray_new:Nn \g__regex_balance_intarray { 65536 }
```

(End definition for `\g__regex_balance_intarray`.)

`__regex_return:` This function triggers either `\prg_return_false:` or `\prg_return_true:` as appropriate to whether a match was found or not. It is used by all user conditionals.

```
26771 \cs_new_protected:Npn \__regex_return:
26772 {
26773   \if_meaning:w \c_true_bool \g__regex_success_bool
26774     \prg_return_true:
26775   \else:
26776     \prg_return_false:
26777   \fi:
26778 }
```

(End definition for `__regex_return:.`)

`__regex_query_set:n` To easily extract subsets of the input once we found the positions at which to cut, store the input tokens one by one into successive `\toks` registers. Also store the brace balance (used to check for overall brace balance) in an array.

```
26779 \cs_new_protected:Npn \__regex_query_set:n #1
26780 {
26781   \int_zero:N \l__regex_balance_int
26782   \int_zero:N \l__regex_curr_pos_int
26783   \__regex_query_set_aux:nN { } F
26784   \tl_analysis_map_inline:nn {#1}
26785     { \__regex_query_set_aux:nN {##1} ##3 }
26786   \__regex_query_set_aux:nN { } F
```

```

26787     \int_set_eq:NN \l__regex_max_pos_int \l__regex_curr_pos_int
26788   }
26789 \cs_new_protected:Npn \__regex_query_set_aux:nN #1#2
26790 {
26791   \int_incr:N \l__regex_curr_pos_int
26792   \__regex_toks_set:Nn \l__regex_curr_pos_int {#1}
26793   \__kernel_intarray_gset:Nnn \g__regex_balance_intarray
26794     { \l__regex_curr_pos_int } { \l__regex_balance_int }
26795   \if_case:w "#2 \exp_stop_f:
26796     \or: \int_incr:N \l__regex_balance_int
26797     \or: \int_decr:N \l__regex_balance_int
26798     \fi:
26799 }

```

(End definition for __regex_query_set:n and __regex_query_set_aux:nN.)

41.7.2 Matching

`__regex_if_match:nn` We don't track submatches, and stop after a single match. Build the NFA with #1, and perform the match on the query #2.

```

26800 \cs_new_protected:Npn \__regex_if_match:nn #1#2
26801 {
26802   \group_begin:
26803     \__regex_disable_submatches:
26804     \__regex_single_match:
26805     #1
26806     \__regex_match:n {#2}
26807   \group_end:
26808 }

```

(End definition for __regex_if_match:nn.)

`__regex_count:nnN` Again, we don't care about submatches. Instead of aborting after the first "longest match" is found, we search for multiple matches, incrementing `\l__regex_match_count_int` every time to record the number of matches. Build the NFA and match. At the end, store the result in the user's variable.

```

26809 \cs_new_protected:Npn \__regex_count:nnN #1#2#3
26810 {
26811   \group_begin:
26812     \__regex_disable_submatches:
26813     \int_zero:N \l__regex_match_count_int
26814     \__regex_multi_match:n { \int_incr:N \l__regex_match_count_int }
26815     #1
26816     \__regex_match:n {#2}
26817     \exp_args:NNNo
26818     \group_end:
26819     \int_set:Nn #3 { \int_use:N \l__regex_match_count_int }
26820 }

```

(End definition for __regex_count:nnN.)

41.7.3 Extracting submatches

`__regex_extract_once:nnN` Match once or multiple times. After each match (or after the only match), extract the submatches using `__regex_extract:.` At the end, store the sequence containing all the submatches into the user variable `#3` after closing the group.

```

26821 \cs_new_protected:Npn __regex_extract_once:nnN #1#2#3
26822 {
26823   \group_begin:
26824     __regex_single_match:
26825     #1
26826     __regex_match:n {#2}
26827     __regex_extract:
26828     __regex_query_set:n {#2}
26829     __regex_group_end_extract_seq:N #3
26830   }
26831 \cs_new_protected:Npn __regex_extract_all:nnN #1#2#3
26832 {
26833   \group_begin:
26834     __regex_multi_match:n { __regex_extract: }
26835     #1
26836     __regex_match:n {#2}
26837     __regex_query_set:n {#2}
26838     __regex_group_end_extract_seq:N #3
26839   }

```

(End definition for `__regex_extract_once:nnN` and `__regex_extract_all:nnN`.)

`__regex_split:nnN` Splitting at submatches is a bit more tricky. For each match, extract all submatches, and replace the zeroth submatch by the part of the query between the start of the match attempt and the start of the zeroth submatch. This is inhibited if the delimiter matched an empty token list at the start of this match attempt. After the last match, store the last part of the token list, which ranges from the start of the match attempt to the end of the query. This step is inhibited if the last match was empty and at the very end: decrement `\l__regex_submatch_int`, which controls which matches will be used.

```

26840 \cs_new_protected:Npn __regex_split:nnN #1#2#3
26841 {
26842   \group_begin:
26843     __regex_multi_match:n
26844     {
26845       \if_int_compare:w
26846         \l__regex_start_pos_int < \l__regex_success_pos_int
26847         __regex_extract:
26848         __kernel_intarray_gset:Nnn \g__regex_submatch_prev_intarray
26849         { \l__regex_zeroth_submatch_int } { 0 }
26850         __kernel_intarray_gset:Nnn \g__regex_submatch_end_intarray
26851         { \l__regex_zeroth_submatch_int }
26852         {
26853           __kernel_intarray_item:Nn \g__regex_submatch_begin_intarray
26854           { \l__regex_zeroth_submatch_int }
26855         }
26856         __kernel_intarray_gset:Nnn \g__regex_submatch_begin_intarray
26857         { \l__regex_zeroth_submatch_int }
26858         { \l__regex_start_pos_int }
26859       \fi:

```

```

26860     }
26861     #1
26862     \__regex_match:n {#2}
26863     \__regex_query_set:n {#2}
26864     \__kernel_intarray_gset:Nnn \g__regex_submatch_prev_intarray
26865     { \l__regex_submatch_int } { 0 }
26866     \__kernel_intarray_gset:Nnn \g__regex_submatch_end_intarray
26867     { \l__regex_submatch_int }
26868     { \l__regex_max_pos_int }
26869     \__kernel_intarray_gset:Nnn \g__regex_submatch_begin_intarray
26870     { \l__regex_submatch_int }
26871     { \l__regex_start_pos_int }
26872     \int_incr:N \l__regex_submatch_int
26873     \if_meaning:w \c_true_bool \l__regex_empty_success_bool
26874     \if_int_compare:w \l__regex_start_pos_int = \l__regex_max_pos_int
26875     \int_decr:N \l__regex_submatch_int
26876     \fi:
26877     \fi:
26878     \__regex_group_end_extract_seq:N #3
26879 }

```

(End definition for __regex_split:nnN.)

__regex_group_end_extract_seq:N The end-points of submatches are stored as entries of two arrays from \l__regex_min_submatch_int to \l__regex_submatch_int (exclusive). Extract the relevant ranges into \l__regex_internal_a_tl. We detect unbalanced results using the two flags __regex_begin and __regex_end, raised whenever we see too many begin-group or end-group tokens in a submatch.

```

26880 \cs_new_protected:Npn \__regex_group_end_extract_seq:N #1
26881 {
26882     \flag_clear:n { __regex_begin }
26883     \flag_clear:n { __regex_end }
26884     \seq_set_from_function:NnN \l__regex_internal_seq
26885     {
26886         \int_step_function:nnN { \l__regex_min_submatch_int }
26887         { \l__regex_submatch_int - 1 }
26888     }
26889     \__regex_extract_seq_aux:n
26890     \int_compare:nNnF
26891     {
26892         \flag_height:n { __regex_begin } +
26893         \flag_height:n { __regex_end }
26894     }
26895     = 0
26896     {
26897         \__kernel_msg_error:nnxxx { kernel } { result-unbalanced }
26898         { splitting~or~extracting~submatches }
26899         { \flag_height:n { __regex_end } }
26900         { \flag_height:n { __regex_begin } }
26901     }
26902     \seq_set_map_x:Nnn \l__regex_internal_seq \l__regex_internal_seq {##1}
26903     \exp_args:NNNo
26904     \group_end:
26905     \tl_set:Nn #1 { \l__regex_internal_seq }

```



```
26906 }
```

(End definition for `_regex_group_end_extract_seq:N`.)

`_regex_extract_seq_aux:n` The `:n` auxiliary builds one item of the sequence of submatches. First compute the brace balance of the submatch, then extract the submatch from the query, adding the appropriate braces and raising a flag if the submatch is not balanced.

```
\_regex_extract_seq_aux:ww
26907 \cs_new:Npn \_regex_extract_seq_aux:n #1
26908 {
26909   \exp_after:wN \_regex_extract_seq_aux:ww
26910   \int_value:w \_regex_submatch_balance:n {#1} ; #1;
26911 }
26912 \cs_new:Npn \_regex_extract_seq_aux:ww #1; #2;
26913 {
26914   \if_int_compare:w #1 < 0 \exp_stop_f:
26915     \flag_raise:n { \_regex_end }
26916     \prg_replicate:nn {-#1} { \exp_not:n { { \if_false: } \fi: } }
26917   \fi:
26918   \_regex_query_submatch:n {#2}
26919   \if_int_compare:w #1 > 0 \exp_stop_f:
26920     \flag_raise:n { \_regex_begin }
26921     \prg_replicate:nn {#1} { \exp_not:n { \if_false: { \fi: } } }
26922   \fi:
26923 }
```

(End definition for `_regex_extract_seq_aux:n` and `_regex_extract_seq_aux:ww`.)

`_regex_extract:` Our task here is to store the list of end-points of submatches, and store them in appropriate array entries, from `\l__regex_zeroth_submatch_int` upwards. First, we store in `\g__regex_submatch_prev_intarray` the position at which the match attempt started. We extract the rest from the comma list `\l__regex_success_submatches_tl`, which starts with entries to be stored in `\g__regex_submatch_begin_intarray` and continues with entries for `\g__regex_submatch_end_intarray`.

```
26924 \cs_new_protected:Npn \_regex_extract:
26925 {
26926   \if_meaning:w \c_true_bool \g__regex_success_bool
26927     \int_set_eq:NN \l__regex_zeroth_submatch_int \l__regex_submatch_int
26928     \prg_replicate:nn \l__regex_capturing_group_int
26929     {
26930       \__kernel_intarray_gset:Nnn \g__regex_submatch_prev_intarray
26931       { \l__regex_submatch_int } { 0 }
26932       \int_incr:N \l__regex_submatch_int
26933     }
26934     \__kernel_intarray_gset:Nnn \g__regex_submatch_prev_intarray
26935     { \l__regex_zeroth_submatch_int } { \l__regex_start_pos_int }
26936     \int_zero:N \l__regex_internal_a_int
26937     \clist_map_inline:Nn \l__regex_success_submatches_tl
26938     {
26939       \if_int_compare:w \l__regex_internal_a_int < \l__regex_capturing_group_int
26940         \__kernel_intarray_gset:Nnn \g__regex_submatch_begin_intarray
26941         { \_regex_int_eval:w \l__regex_zeroth_submatch_int + \l__regex_internal_a_int }
26942       \else:
26943         \__kernel_intarray_gset:Nnn \g__regex_submatch_end_intarray
26944         { \_regex_int_eval:w \l__regex_zeroth_submatch_int + \l__regex_internal_a_int }
```

```

26945         \fi:
26946         \int_incr:N \l__regex_internal_a_int
26947     }
26948     \fi:
26949 }

```

(End definition for __regex_extract:.)

41.7.4 Replacement

__regex_replace_once:nnN Build the NFA and the replacement functions, then find a single match. If the match failed, simply exit the group. Otherwise, we do the replacement. Extract submatches. Compute the brace balance corresponding to replacing this match by the replacement (this depends on submatches). Prepare the replaced token list: the replacement function produces the tokens from the start of the query to the start of the match and the replacement text for this match; we need to add the tokens from the end of the match to the end of the query. Finally, store the result in the user's variable after closing the group: this step involves an additional x-expansion, and checks that braces are balanced in the final result.

```

26950 \cs_new_protected:Npn \__regex_replace_once:nnN #1#2#3
26951 {
26952     \group_begin:
26953     \__regex_single_match:
26954     #1
26955     \exp_args:No \__regex_match:n {#3}
26956     \if_meaning:w \c_false_bool \g__regex_success_bool
26957     \group_end:
26958     \else:
26959     \__regex_extract:
26960     \exp_args:No \__regex_query_set:n {#3}
26961     \__regex_replacement:n {#2}
26962     \int_set:Nn \l__regex_balance_int
26963     {
26964         \__regex_replacement_balance_one_match:n
26965         { \l__regex_zeroth_submatch_int }
26966     }
26967     \__kernel_tl_set:Nx \l__regex_internal_a_tl
26968     {
26969         \__regex_replacement_do_one_match:n
26970         { \l__regex_zeroth_submatch_int }
26971         \__regex_query_range:nn
26972         {
26973             \__kernel_intarray_item:Nn \g__regex_submatch_end_intarray
26974             { \l__regex_zeroth_submatch_int }
26975         }
26976         { \l__regex_max_pos_int }
26977     }
26978     \__regex_group_end_replace:N #3
26979     \fi:
26980 }

```

(End definition for __regex_replace_once:nnN.)

__regex_replace_all:nnN Match multiple times, and for every match, extract submatches and additionally store the position at which the match attempt started. The entries from \l__regex_min_submatch_int to \l__regex_submatch_int hold information about submatches of every

match in order; each match corresponds to `\l__regex_capturing_group_int` consecutive entries. Compute the brace balance corresponding to doing all the replacements: this is the sum of brace balances for replacing each match. Join together the replacement texts for each match (including the part of the query before the match), and the end of the query.

```

26981 \cs_new_protected:Npn \__regex_replace_all:nnN #1#2#3
26982 {
26983   \group_begin:
26984   \__regex_multi_match:n { \__regex_extract: }
26985   #1
26986   \exp_args:No \__regex_match:n {#3}
26987   \exp_args:No \__regex_query_set:n {#3}
26988   \__regex_replacement:n {#2}
26989   \int_set:Nn \l__regex_balance_int
26990   {
26991     0
26992     \int_step_function:nnnN
26993     { \l__regex_min_submatch_int }
26994     \l__regex_capturing_group_int
26995     { \l__regex_submatch_int - 1 }
26996     \__regex_replacement_balance_one_match:n
26997   }
26998   \__kernel_tl_set:Nx \l__regex_internal_a_tl
26999   {
27000     \int_step_function:nnnN
27001     { \l__regex_min_submatch_int }
27002     \l__regex_capturing_group_int
27003     { \l__regex_submatch_int - 1 }
27004     \__regex_replacement_do_one_match:n
27005     \__regex_query_range:nn
27006     \l__regex_start_pos_int \l__regex_max_pos_int
27007   }
27008   \__regex_group_end_replace:N #3
27009 }

```

(End definition for `__regex_replace_all:nnN`.)

`__regex_group_end_replace:N` If the brace balance is not 0, raise an error. Then set the user's variable `#1` to the x-expansion of `\l__regex_internal_a_tl`, adding the appropriate braces to produce a balanced result. And end the group.

```

27010 \cs_new_protected:Npn \__regex_group_end_replace:N #1
27011 {
27012   \if_int_compare:w \l__regex_balance_int = 0 \exp_stop_f:
27013   \else:
27014     \__kernel_msg_error:nnxxx { kernel } { result-unbalanced }
27015     { replacing }
27016     { \int_max:nn { - \l__regex_balance_int } { 0 } }
27017     { \int_max:nn { \l__regex_balance_int } { 0 } }
27018   \fi:
27019   \use:x
27020   {
27021     \group_end:
27022     \tl_set:Nn \exp_not:N #1
27023     {

```

```

27024         \if_int_compare:w \l__regex_balance_int < 0 \exp_stop_f:
27025         \prg_replicate:nn { - \l__regex_balance_int }
27026         { { \if_false: } \fi: }
27027     \fi:
27028     \l__regex_internal_a_tl
27029     \if_int_compare:w \l__regex_balance_int > 0 \exp_stop_f:
27030     \prg_replicate:nn { \l__regex_balance_int }
27031     { \if_false: { \fi: } }
27032     \fi:
27033 }
27034 }
27035 }

```

(End definition for `__regex_group_end_replace:N`.)

41.7.5 Peeking ahead

`\l__regex_peek_true_tl` True/false code arguments of `\peek_regex:nTF` or similar.

```

\l__regex_peek_false_tl 27036 \tl_new:N \l__regex_peek_true_tl
27037 \tl_new:N \l__regex_peek_false_tl

```

(End definition for `\l__regex_peek_true_tl` and `\l__regex_peek_false_tl`.)

`\l__regex_replacement_tl` When peeking in `\peek_regex_replace_once:nnTF` we need to store the replacement text.

```
27038 \tl_new:N \l__regex_replacement_tl
```

(End definition for `\l__regex_replacement_tl`.)

`\l__regex_input_tl` Stores each token found as `__regex_input_item:n` $\{\langle tokens \rangle\}$, where the $\langle tokens \rangle$ o-expand to the token found, as for `\tl_analysis_map_inline:nn`.

```

\__regex_input_item:n 27039 \tl_new:N \l__regex_input_tl
27040 \cs_new_eq:NN \__regex_input_item:n ?

```

(End definition for `\l__regex_input_tl` and `__regex_input_item:n`.)

`\peek_regex:nTF`

`\peek_regex:NTF`

`\peek_regex_remove_once:nTF`

`\peek_regex_remove_once:NTF`

The T and F functions just call the corresponding TF function. The four TF functions differ along two axes: whether to remove the token or not, distinguished by using `__regex_peek_end:` or `__regex_peek_remove_end:n` (the latter case needs an argument, as we will see), and whether the regex has to be compiled or is already in an N-type variable, distinguished by calling `__regex_build_aux:Nn` or `__regex_build_aux:NN`. The first argument of these functions is `\c_false_bool` to indicate that there should be no implicit insertion of a wildcard at the start of the pattern: otherwise the code would keep looking further into the input stream until matching the regex.

```

27041 \cs_new_protected:Npn \peek_regex:nTF #1
27042 {
27043     \__regex_peek:nnTF
27044     { \__regex_build_aux:Nn \c_false_bool {#1} }
27045     { \__regex_peek_end: }
27046 }
27047 \cs_new_protected:Npn \peek_regex:nT #1#2
27048 { \peek_regex:nTF {#1} {#2} { } }
27049 \cs_new_protected:Npn \peek_regex:nF #1 { \peek_regex:nTF {#1} { } }
27050 \cs_new_protected:Npn \peek_regex:NTF #1

```

```

27051 {
27052     \__regex_peek:nnTF
27053     { \__regex_build_aux:NN \c_false_bool #1 }
27054     { \__regex_peek_end: }
27055 }
27056 \cs_new_protected:Npn \peek_regex:NT #1#2
27057 { \peek_regex:NTF #1 {#2} { } }
27058 \cs_new_protected:Npn \peek_regex:NF #1 { \peek_regex:NTF {#1} { } }
27059 \cs_new_protected:Npn \peek_regex_remove_once:nTF #1
27060 {
27061     \__regex_peek:nnTF
27062     { \__regex_build_aux:Nn \c_false_bool {#1} }
27063     { \__regex_peek_remove_end:n {##1} }
27064 }
27065 \cs_new_protected:Npn \peek_regex_remove_once:nT #1#2
27066 { \peek_regex_remove_once:NTF {#1} {#2} { } }
27067 \cs_new_protected:Npn \peek_regex_remove_once:nF #1
27068 { \peek_regex_remove_once:NTF {#1} { } }
27069 \cs_new_protected:Npn \peek_regex_remove_once:NTF #1
27070 {
27071     \__regex_peek:nnTF
27072     { \__regex_build_aux:NN \c_false_bool #1 }
27073     { \__regex_peek_remove_end:n {##1} }
27074 }
27075 \cs_new_protected:Npn \peek_regex_remove_once:NT #1#2
27076 { \peek_regex_remove_once:NTF #1 {#2} { } }
27077 \cs_new_protected:Npn \peek_regex_remove_once:NF #1
27078 { \peek_regex_remove_once:NTF #1 { } }

```

(End definition for `\peek_regex:nTF` and others. These functions are documented on page 142.)

`__regex_peek:nnTF` Store the user's true/false codes (plus `\group_end:`) into two token lists. Then build the automaton with #1, without submatch tracking, and aiming for a single match. Then start matching by setting up a few variables like for any regex matching like `\regex-match:nnTF`, with the addition of `\l__regex_input_tl` that keeps track of the tokens seen, to reinsert them at the end. Instead of `\tl_analysis_map_inline:nn` on the input, we call `\peek_analysis_map_inline:n` to go through tokens in the input stream. Since `__regex_match_one_token:nnN` calls `__regex_maplike_break:` we need to catch that and break the `\peek_analysis_map_inline:n` loop instead.

```

27079 \cs_new_protected:Npn \__regex_peek:nnTF #1
27080 {
27081     \__regex_peek_aux:nnTF
27082     {
27083         \__regex_disable_submatches:
27084         #1
27085     }
27086 }
27087 \cs_new_protected:Npn \__regex_peek_aux:nnTF #1#2#3#4
27088 {
27089     \group_begin:
27090     \tl_set:Nn \l__regex_peek_true_tl { \group_end: #3 }
27091     \tl_set:Nn \l__regex_peek_false_tl { \group_end: #4 }
27092     \__regex_single_match:
27093     #1

```

```

27094     \__regex_match_init:
27095     \tl_build_clear:N \l__regex_input_tl
27096     \__regex_match_once_init:
27097     \peek_analysis_map_inline:n
27098     {
27099         \tl_build_put_right:Nn \l__regex_input_tl
27100         { \__regex_input_item:n {##1} }
27101         \__regex_match_one_token:nnN {##1} {##2} ##3
27102         \use_none:nnn
27103         \prg_break_point:Nn \__regex_maplike_break:
27104         { \peek_analysis_map_break:n {#2} }
27105     }
27106 }

```

(End definition for __regex_peek:nnTF and __regex_peek_aux:nnTF.)

__regex_peek_end: Once the regex matches (or permanently fails to match) we call __regex_peek_end:, or
 __regex_peek_remove_end:n __regex_peek_remove_end:n with argument the last token seen. For \peek_regex:nTF
 we reinsert tokens seen by calling __regex_peek_reinsert:N regardless of the result
 of the match. For \peek_regex_remove_once:nTF we reinsert the tokens seen only if
 the match failed; otherwise we just reinsert the tokens #1, with one expansion. To be
 more precise, #1 consists of tokens that o-expand and x-expand to the last token seen,
 for example it is \exp_not:N <cs> for a control sequence. This means that just doing
 \exp_after:wN \l__regex_peek_true_tl #1 would be unsafe because the expansion of
 <cs> would be suppressed.

```

27107 \cs_new_protected:Npn \__regex_peek_end:
27108 {
27109     \bool_if:NTF \g__regex_success_bool
27110     { \__regex_peek_reinsert:N \l__regex_peek_true_tl }
27111     { \__regex_peek_reinsert:N \l__regex_peek_false_tl }
27112 }
27113 \cs_new_protected:Npn \__regex_peek_remove_end:n #1
27114 {
27115     \bool_if:NTF \g__regex_success_bool
27116     { \exp_args:NNo \use:nn \l__regex_peek_true_tl {#1} }
27117     { \__regex_peek_reinsert:N \l__regex_peek_false_tl }
27118 }

```

(End definition for __regex_peek_end: and __regex_peek_remove_end:n.)

__regex_peek_reinsert:N Insert the true/false code #1, followed by the tokens found, which were stored in \l__-
 __regex_reinsert_item:n regex_input_tl. For this, loop through that token list using __regex_reinsert_-
 item:n, which expands #1 once to get a single token, and jumps over it to expand
 what follows, with suitable \exp:w and \exp_end:. We cannot just use \use:e on the
 whole token list because the result may be unbalanced, which would stop the primitive
 prematurely, or let it continue beyond where we would like.

```

27119 \cs_new_protected:Npn \__regex_peek_reinsert:N #1
27120 {
27121     \tl_build_end:N \l__regex_input_tl
27122     \cs_set_eq:NN \__regex_input_item:n \__regex_reinsert_item:n
27123     \exp_after:wN #1 \exp:w \l__regex_input_tl \exp_end:
27124 }
27125 \cs_new_protected:Npn \__regex_reinsert_item:n #1

```

```

27126 {
27127     \exp_after:wN \exp_after:wN
27128     \exp_after:wN \exp_end:
27129     \exp_after:wN \exp_after:wN
27130     #1
27131     \exp:w
27132 }

```

(End definition for `_regex_peek_reinsert:N` and `_regex_reinsert_item:n`.)

`\peek_regex_replace_once:nn`
`\peek_regex_replace_once:nnTF`
`\peek_regex_replace_once:Nn`
`\peek_regex_replace_once:NnTF`

Similar to `\peek_regex:nTF` above.

```

27133 \cs_new_protected:Npn \peek_regex_replace_once:nnTF #1
27134 { \_regex_peek_replace:nnTF { \_regex_build_aux:Nn \c_false_bool {#1} } }
27135 \cs_new_protected:Npn \peek_regex_replace_once:nnT #1#2#3
27136 { \peek_regex_replace_once:nnTF {#1} {#2} {#3} { } }
27137 \cs_new_protected:Npn \peek_regex_replace_once:nnF #1#2
27138 { \peek_regex_replace_once:nnTF {#1} {#2} { } }
27139 \cs_new_protected:Npn \peek_regex_replace_once:nn #1#2
27140 { \peek_regex_replace_once:nnTF {#1} {#2} { } { } }
27141 \cs_new_protected:Npn \peek_regex_replace_once:NnTF #1
27142 { \_regex_peek_replace:nnTF { \_regex_build_aux:Nn \c_false_bool #1 } }
27143 \cs_new_protected:Npn \peek_regex_replace_once:NnT #1#2#3
27144 { \peek_regex_replace_once:NnTF #1 {#2} {#3} { } }
27145 \cs_new_protected:Npn \peek_regex_replace_once:NnF #1#2
27146 { \peek_regex_replace_once:NnTF #1 {#2} { } }
27147 \cs_new_protected:Npn \peek_regex_replace_once:Nn #1#2
27148 { \peek_regex_replace_once:NnTF #1 {#2} { } { } }

```

(End definition for `\peek_regex_replace_once:nnTF` and `\peek_regex_replace_once:NnTF`. These functions are documented on page 143.)

`_regex_peek_replace:nnTF`

Same as `_regex_peek:nnTF` (used for `\peek_regex:nTF` above), but without disabling submatches, and with a different end. The replacement text #2 is stored, to be analyzed later.

```

27149 \cs_new_protected:Npn \_regex_peek_replace:nnTF #1#2
27150 {
27151     \tl_set:Nn \l__regex_replacement_tl {#2}
27152     \_regex_peek_aux:nnTF {#1} { \_regex_peek_replace_end: }
27153 }

```

(End definition for `_regex_peek_replace:nnTF`.)

`_regex_peek_replace_end:`

If the match failed `_regex_peek_reinsert:N` reinserts the tokens found. Otherwise, finish storing the submatch information using `_regex_extract:`, and store the input into `\toks`. Redefine a few auxiliaries to change slightly their expansion behaviour as explained below. Analyse the replacement text with `_regex_replacement:n`, which as usual defines `_regex_replacement_do_one_match:n` to insert the tokens from the start of the match attempt to the beginning of the match, followed by the replacement text. The `\use:x` expands for instance the trailing `_regex_query_range:nn` down to a sequence of `_regex_reinsert_item:n` $\{\langle tokens \rangle\}$ where $\langle tokens \rangle$ o-expand to a single token that we want to insert. After x-expansion, `\use:x` does `\use:n`, so we have `\exp_after:wN \l__regex_peek_true_tl \exp:w ... \exp_end:.` This is set up such as to obtain `\l__regex_peek_true_tl` followed by the replaced tokens (possibly unbalanced) in the input stream.

```

27154 \cs_new_protected:Npn \__regex_peek_replace_end:
27155 {
27156   \bool_if:NTF \g__regex_success_bool
27157   {
27158     \__regex_extract:
27159     \__regex_query_set_from_input_tl:
27160     \cs_set_eq:NN \__regex_replacement_put:n \__regex_peek_replacement_put:n
27161     \cs_set_eq:NN \__regex_replacement_put_submatch_aux:n
27162     \__regex_peek_replacement_put_submatch_aux:n
27163     \cs_set_eq:NN \__regex_input_item:n \__regex_reinsert_item:n
27164     \cs_set_eq:NN \__regex_replacement_exp_not:N \__regex_peek_replacement_token:n
27165     \cs_set_eq:NN \__regex_replacement_exp_not:V \__regex_peek_replacement_var:N
27166     \exp_args:No \__regex_replacement:n { \l__regex_replacement_tl }
27167     \use:x
27168     {
27169       \exp_not:n { \exp_after:wN \l__regex_peek_true_tl \exp:w }
27170       \__regex_replacement_do_one_match:n
27171       { \l__regex_zeroth_submatch_int }
27172       \__regex_query_range:nn
27173       {
27174         \__kernel_intarray_item:Nn \g__regex_submatch_end_intarray
27175         { \l__regex_zeroth_submatch_int }
27176       }
27177       { \l__regex_max_pos_int }
27178     }
27179   }
27180 }
27181 { \__regex_peek_reinsert:N \l__regex_peek_false_tl }
27182 }

```

(End definition for __regex_peek_replace_end:.)

__regex_query_set_from_input_tl: The input was stored into \l__regex_input_tl as successive items __regex_input_item:n {<tokens>}. Store that in successive \toks. It's not clear whether the empty entries before and after are both useful.

```

27183 \cs_new_protected:Npn \__regex_query_set_from_input_tl:
27184 {
27185   \tl_build_end:N \l__regex_input_tl
27186   \int_zero:N \l__regex_curr_pos_int
27187   \cs_set_eq:NN \__regex_input_item:n \__regex_query_set_item:n
27188   \__regex_query_set_item:n { }
27189   \l__regex_input_tl
27190   \__regex_query_set_item:n { }
27191   \int_set_eq:NN \l__regex_max_pos_int \l__regex_curr_pos_int
27192 }
27193 \cs_new_protected:Npn \__regex_query_set_item:n #1
27194 {
27195   \int_incr:N \l__regex_curr_pos_int
27196   \__regex_toks_set:Nn \l__regex_curr_pos_int { \__regex_input_item:n {#1} }
27197 }

```

(End definition for __regex_query_set_from_input_tl: and __regex_query_set_item:n.)

__regex_peek_replacement_put:n While building the replacement function __regex_replacement_do_one_match:n, we often want to put simple material, given as #1, whose x-expansion o-expands to a single

token. Normally we can just add the token to `\l__regex_build_tl`, but for `\peek_regex_replace_once:nnTF` we eventually want to do some strange expansion that is basically using `\exp_after:wN` to jump through numerous tokens (we cannot use `x`-expansion like for `\regex_replace_once:nnNTF` because it is ok for the result to be unbalanced since we insert it in the input stream rather than storing it. When within a csname we don't do any such shenanigan because `\cs:w ... \cs_end:` does all the expansion we need.

```

27198 \cs_new_protected:Npn \__regex_peek_replacement_put:n #1
27199 {
27200   \if_case:w \l__regex_replacement_csnames_int
27201     \tl_build_put_right:Nn \l__regex_build_tl
27202     { \exp_not:N \__regex_reinsert_item:n {#1} }
27203   \else:
27204     \tl_build_put_right:Nn \l__regex_build_tl {#1}
27205   \fi:
27206 }

```

(End definition for `__regex_peek_replacement_put:n`.)

`__regex_peek_replacement_token:n` When hit with `\exp:w, __regex_peek_replacement_token:n {<token>}` stops `\exp_end:` and does `\exp_after:wN <token> \exp:w` to continue expansion after it.

```

27207 \cs_new_protected:Npn \__regex_peek_replacement_token:n #1
27208 { \exp_after:wN \exp_end: \exp_after:wN #1 \exp:w }

```

(End definition for `__regex_peek_replacement_token:n`.)

`__regex_peek_replacement_put_submatch_aux:n` While analyzing the replacement we also have to insert submatches found in the query. Since query items `__regex_input_item:n {<tokens>}` expand correctly only when surrounded by `\exp:w ... \exp_end:`, and since these expansion controls are not there within csnames (because `\cs:w ... \cs_end:` make them unnecessary in most cases), we have to put `\exp:w` and `\exp_end:` by hand here.

```

27209 \cs_new_protected:Npn \__regex_peek_replacement_put_submatch_aux:n #1
27210 {
27211   \if_case:w \l__regex_replacement_csnames_int
27212     \tl_build_put_right:Nn \l__regex_build_tl
27213     { \__regex_query_submatch:n { \int_eval:n { #1 + ##1 } } }
27214   \else:
27215     \tl_build_put_right:Nn \l__regex_build_tl
27216     { \exp:w \__regex_query_submatch:n { \int_eval:n { #1 + ##1 } } \exp_end: }
27217   \fi:
27218 }

```

(End definition for `__regex_peek_replacement_put_submatch_aux:n`.)

`__regex_peek_replacement_var:N` This is used for `\u` outside csnames. It makes sure to continue expansion with `\exp:w` before expanding the variable `#1` and stopping the `\exp:w` that precedes.

```

27219 \cs_new_protected:Npn \__regex_peek_replacement_var:N #1
27220 {
27221   \exp_after:wN \exp_last_unbraced:NV
27222   \exp_after:wN \exp_end:
27223   \exp_after:wN #1
27224   \exp:w
27225 }

```

(End definition for `__regex_peek_replacement_var:N`.)

41.8 Messages

Messages for the preparsing phase.

```
27226 \use:x
27227 {
27228   \__kernel_msg_new:nnn { kernel } { trailing-backslash }
27229   { Trailing-escape-char~'\iow_char:N\\'~in-regex-or-replacement. }
27230   \__kernel_msg_new:nnn { kernel } { x-missing-rbrace }
27231   {
27232     Missing~brace~'\iow_char:N\}'~in-regex~
27233     '...\iow_char:N\x\iow_char:N\{...#1'.
27234   }
27235   \__kernel_msg_new:nnn { kernel } { x-overflow }
27236   {
27237     Character~code~##1~too~large~in~
27238     \iow_char:N\x\iow_char:N\{##2\iow_char:N\}~regex.
27239   }
27240 }
```

Invalid quantifier.

```
27241 \__kernel_msg_new:nnnn { kernel } { invalid-quantifier }
27242 { Braced~quantifier~'~#1'~may~not~be~followed~by~'~#2'. }
27243 {
27244   The~character~'~#2'~is~invalid~in~the~braced~quantifier~'~#1'.~
27245   The~only~valid~quantifiers~are~'*',~'?','+',~'<int>',~
27246   '~<min>',~'~and~'~<min>,<max>',~optionally~followed~by~'?''.
27247 }
```

Messages for missing or extra closing brackets and parentheses, with some fancy singular/plural handling for the case of parentheses.

```
27248 \__kernel_msg_new:nnnn { kernel } { missing-rbrack }
27249 { Missing~right~bracket~inserted~in~regular-expression. }
27250 {
27251   LaTeX~was~given~a~regular~expression~where~a~character~class~
27252   was~started~with~'~['~,~but~the~matching~'~']'~is~missing.
27253 }
27254 \__kernel_msg_new:nnnn { kernel } { missing-rparen }
27255 {
27256   Missing~right~
27257   \int_compare:nTF { #1 = 1 } { parenthesis } { parentheses } ~
27258   inserted~in~regular-expression.
27259 }
27260 {
27261   LaTeX~was~given~a~regular~expression~with~\int_eval:n {#1} ~
27262   more~left~parentheses~than~right~parentheses.
27263 }
27264 \__kernel_msg_new:nnnn { kernel } { extra-rparen }
27265 { Extra~right~parenthesis~ignored~in~regular-expression. }
27266 {
27267   LaTeX~came~across~a~closing~parenthesis~when~no~submatch~group~
27268   was~open.~The~parenthesis~will~be~ignored.
27269 }
```

Some escaped alphanumerics are not allowed everywhere.

```
27270 \__kernel_msg_new:nnnn { kernel } { bad-escape }
```

```

27271 {
27272     Invalid-escape~'\iow_char:N\\#1'~
27273     \_regex_if_in_cs:TF { within-a-control~sequence. }
27274     {
27275         \_regex_if_in_class:TF
27276         { in-a-character-class. }
27277         { following-a-category~test. }
27278     }
27279 }
27280 {
27281     The-escape~sequence~'\iow_char:N\\#1'~may-not-appear~
27282     \_regex_if_in_cs:TF
27283     {
27284         within-a-control~sequence~test~introduced-by~
27285         '\iow_char:N\\c\iow_char:N\{'~
27286     }
27287     {
27288         \_regex_if_in_class:TF
27289         { within-a-character-class~
27290         { following-a-category~test~such-as~'\iow_char:N\\cL'~ }
27291         because-it~does~not~match~exactly~one~character.
27292     }
27293 }

```

Range errors.

```

27294 \_kernel_msg_new:nnnn { kernel } { range-missing-end }
27295 { Invalid-end-point~for~range~'#1-#2'~in~character~class. }
27296 {
27297     The-end-point~'#2'~of~the~range~'#1-#2'~may-not-serve-as-an~
27298     end-point~for~a~range:-alphanumeric~characters-should-not-be~
27299     escaped,~and-non-alphanumeric~characters~should-be-escaped.
27300 }
27301 \_kernel_msg_new:nnnn { kernel } { range-backwards }
27302 { Range~'[#1-#2]'~out-of-order~in~character~class. }
27303 {
27304     In-ranges~of~characters~'[x-y]'~appearing-in~character~classes,~
27305     the~first~character~code~must~not~be~larger~than~the~second.~
27306     Here,~'#1'~has~character~code~\int_eval:n {'#1},~while~
27307     '#2'~has~character~code~\int_eval:n {'#2}.
27308 }

```

Errors related to \c and \u.

```

27309 \_kernel_msg_new:nnnn { kernel } { c-bad-mode }
27310 { Invalid-nested~'\iow_char:N\\c'~escape~in~regular~expression. }
27311 {
27312     The~'\iow_char:N\\c'~escape~cannot~be~used~within~
27313     a~control~sequence~test~'\iow_char:N\\c{...}'~
27314     nor~another~category~test.~
27315     To~combine~several~category~tests,~use~'\iow_char:N\\c[...]'~
27316 }
27317 \_kernel_msg_new:nnnn { kernel } { c-C-invalid }
27318 { '\iow_char:N\\c'~should-be-followed-by~'.'~or~'(',~not~'#1'. }
27319 {
27320     The~'\iow_char:N\\cC'~construction~restricts~the~next~item~to~be~a~
27321     control~sequence~or~the~next~group~to~be~made~of~control~sequences.~

```

```

27322     It~only~makes~sense~to~follow~it~by~'.'~or~by~a~group.
27323 }
27324 \__kernel_msg_new:nnnn { kernel } { c-lparen-in-class }
27325 { Catcode~test~cannot~apply~to~group~in~character~class }
27326 {
27327     Construction~such~as~'\iow_char:N\\cL(abc)''~are~not~allowed~inside~a~
27328     class~'[...]''~because~classes~do~not~match~multiple~characters~at~once.
27329 }
27330 \__kernel_msg_new:nnnn { kernel } { c-missing-rbrace }
27331 { Missing~right~brace~inserted~for~'\iow_char:N\\c''~escape. }
27332 {
27333     LaTeX~was~given~a~regular~expression~where~a~
27334     '\iow_char:N\\c\iow_char:N\{...'~construction~was~not~ended~
27335     with~a~closing~brace~'\iow_char:N\}''.
27336 }
27337 \__kernel_msg_new:nnnn { kernel } { c-missing-rbrack }
27338 { Missing~right~bracket~inserted~for~'\iow_char:N\\c''~escape. }
27339 {
27340     A~construction~'\iow_char:N\\c[...'~appears~in~a~
27341     regular~expression,~but~the~closing~'~'~is~not~present.
27342 }
27343 \__kernel_msg_new:nnnn { kernel } { c-missing-category }
27344 { Invalid~character~'#1''~following~'\iow_char:N\\c''~escape. }
27345 {
27346     In~regular~expressions,~the~'\iow_char:N\\c''~escape~sequence~
27347     may~only~be~followed~by~a~left~brace,~a~left~bracket,~or~a~
27348     capital~letter~representing~a~character~category,~namely~
27349     one~of~'ABCDELMOPTU'.
27350 }
27351 \__kernel_msg_new:nnnn { kernel } { c-trailing }
27352 { Trailing~category~code~escape~'\iow_char:N\\c'... }
27353 {
27354     A~regular~expression~ends~with~'\iow_char:N\\c''~followed~
27355     by~a~letter.~It~will~be~ignored.
27356 }
27357 \__kernel_msg_new:nnnn { kernel } { u-missing-lbrace }
27358 { Missing~left~brace~following~'\iow_char:N\\u''~escape. }
27359 {
27360     The~'\iow_char:N\\u''~escape~sequence~must~be~followed~by~
27361     a~brace~group~with~the~name~of~the~variable~to~use.
27362 }
27363 \__kernel_msg_new:nnnn { kernel } { u-missing-rbrace }
27364 { Missing~right~brace~inserted~for~'\iow_char:N\\u''~escape. }
27365 {
27366     LaTeX~
27367     \str_if_eq:eeTF { } {#2}
27368     { reached~the~end~of~the~string~ }
27369     { encountered~an~escaped~alphanumeric~character '\iow_char:N\\#2'~ }
27370     when~parsing~the~argument~of~an~
27371     '\iow_char:N\\u\iow_char:N\{...}'~escape.
27372 }

```

Errors when encountering the POSIX syntax [:...:].

```

27373 \__kernel_msg_new:nnnn { kernel } { posix-unsupported }
27374 { POSIX~collating~element~'[#1 ~ #1]''~not~supported. }

```

```

27375 {
27376   The~'[.foo.]'~and~'[=bar=]'~syntaxes~have~a~special~meaning~
27377   in~POSIX~regular~expressions.~This~is~not~supported~by~LaTeX.~
27378   Maybe~you~forgot~to~escape~a~left~bracket~in~a~character~class?
27379 }
27380 \__kernel_msg_new:nnnn { kernel } { posix-unknown }
27381 { POSIX~class~'[:#1:]'~unknown. }
27382 {
27383   '[:#1:]'~is~not~among~the~known~POSIX~classes~
27384   '[:alnum:]',~'[:alpha:]',~'[:ascii:]',~'[:blank:]',~
27385   '[:cntrl:]',~'[:digit:]',~'[:graph:]',~'[:lower:]',~
27386   '[:print:]',~'[:punct:]',~'[:space:]',~'[:upper:]',~
27387   '[:word:]',~and~'[:xdigit:]'.
27388 }
27389 \__kernel_msg_new:nnnn { kernel } { posix-missing-close }
27390 { Missing~closing~':'~for~POSIX~class. }
27391 { The~POSIX~syntax~'#1'~must~be~followed~by~':'~',~not~'#2'. }

```

In various cases, the result of a `l3regex` operation can leave us with an unbalanced token list, which we must re-balance by adding begin-group or end-group character tokens.

```

27392 \__kernel_msg_new:nnnn { kernel } { result-unbalanced }
27393 { Missing~brace~inserted~when~#1. }
27394 {
27395   LaTeX~was~asked~to~do~some~regular~expression~operation,~
27396   and~the~resulting~token~list~would~not~have~the~same~number~
27397   of~begin~group~and~end~group~tokens.~Braces~were~inserted:~
27398   #2~left,~#3~right.
27399 }

```

Error message for unknown options.

```

27400 \__kernel_msg_new:nnnn { kernel } { unknown-option }
27401 { Unknown~option~'#1'~for~regular~expressions. }
27402 {
27403   The~only~available~option~is~'case-insensitive',~toggled~by~
27404   '(?i)'~and~'(?-i)'.
27405 }
27406 \__kernel_msg_new:nnnn { kernel } { special-group-unknown }
27407 { Unknown~special~group~'#1~...'~in~a~regular~expression. }
27408 {
27409   The~only~valid~constructions~starting~with~'(?~are~
27410   '(:~...'~)',~'(?|~...'~)',~'(?i)',~and~'(?-i)'.
27411 }

```

Errors in the replacement text.

```

27412 \__kernel_msg_new:nnnn { kernel } { replacement-c }
27413 { Misused~'\iow_char:N\c'~command~in~a~replacement~text. }
27414 {
27415   In~a~replacement~text,~the~'\iow_char:N\c'~escape~sequence~
27416   can~be~followed~by~one~of~the~letters~'ABCDELMOPTU'~
27417   or~a~brace~group,~not~by~'#1'.
27418 }
27419 \__kernel_msg_new:nnnn { kernel } { replacement-u }
27420 { Misused~'\iow_char:N\u'~command~in~a~replacement~text. }
27421 {

```

```

27422     In~a~replacement~text,~the~'\iow_char:N\\u'~escape~sequence~
27423     must~be~followed~by~a~brace~group~holding~the~name~of~the~
27424     variable~to~use.
27425   }
27426   \__kernel_msg_new:nnnn { kernel } { replacement-g }
27427   {
27428     Missing~brace~for~the~'\iow_char:N\\g'~construction~
27429     in~a~replacement~text.
27430   }
27431   {
27432     In~the~replacement~text~for~a~regular~expression~search,~
27433     submatches~are~represented~either~as~'\iow_char:N \\g{dd..d}',~
27434     or~'\d',~where~'d'~are~single~digits.~Here,~a~brace~is~missing.
27435   }
27436   \__kernel_msg_new:nnnn { kernel } { replacement-catcode-end }
27437   {
27438     Missing~character~for~the~'\iow_char:N\\c<category><character>'~
27439     construction~in~a~replacement~text.
27440   }
27441   {
27442     In~a~replacement~text,~the~'\iow_char:N\\c'~escape~sequence~
27443     can~be~followed~by~one~of~the~letters~'ABCDELMOPTU'~representing~
27444     the~character~category.~Then,~a~character~must~follow.~LaTeX~
27445     reached~the~end~of~the~replacement~when~looking~for~that.
27446   }
27447   \__kernel_msg_new:nnnn { kernel } { replacement-catcode-escaped }
27448   {
27449     Escaped~letter~or~digit~after~category~code~in~replacement~text.
27450   }
27451   {
27452     In~a~replacement~text,~the~'\iow_char:N\\c'~escape~sequence~
27453     can~be~followed~by~one~of~the~letters~'ABCDELMOPTU'~representing~
27454     the~character~category.~Then,~a~character~must~follow,~not~
27455     '\iow_char:N\\#2'.
27456   }
27457   \__kernel_msg_new:nnnn { kernel } { replacement-catcode-in-cs }
27458   {
27459     Category~code~'\iow_char:N\\c#1#3'~ignored~inside~
27460     '\iow_char:N\\c\{...\}'~in~a~replacement~text.
27461   }
27462   {
27463     In~a~replacement~text,~the~category~codes~of~the~argument~of~
27464     '\iow_char:N\\c\{...\}'~are~ignored~when~building~the~control~
27465     sequence~name.
27466   }
27467   \__kernel_msg_new:nnnn { kernel } { replacement-null-space }
27468   { TeX~cannot~build~a~space~token~with~character~code~0. }
27469   {
27470     You~asked~for~a~character~token~with~category~space,~
27471     and~character~code~0,~for~instance~through~
27472     '\iow_char:N\\cS\iow_char:N\\x00'.~
27473     This~specific~case~is~impossible~and~will~be~replaced~
27474     by~a~normal~space.
27475   }

```

```

27476 \__kernel_msg_new:nnnn { kernel } { replacement-missing-rbrace }
27477 { Missing~right~brace-inserted-in-replacement-text. }
27478 {
27479     There~ \int_compare:nTF { #1 = 1 } { was } { were } ~ #1~
27480     missing-right~\int_compare:nTF { #1 = 1 } { brace } { braces } .
27481 }
27482 \__kernel_msg_new:nnnn { kernel } { replacement-missing-rparen }
27483 { Missing~right~parenthesis-inserted-in-replacement-text. }
27484 {
27485     There~ \int_compare:nTF { #1 = 1 } { was } { were } ~ #1~
27486     missing-right~
27487     \int_compare:nTF { #1 = 1 } { parenthesis } { parentheses } .
27488 }

```

Used when showing a regex.

```

27489 \__kernel_msg_new:nnn { kernel } { show-regex }
27490 {
27491     >~Compiled~regex~
27492     \tl_if_empty:nTF {#1} { variable~ #2 } { {#1} } :
27493     #3
27494 }

```

__regex_msg_repeated:nnN This is not technically a message, but seems related enough to go there. The arguments are: #1 is the minimum number of repetitions; #2 is the number of allowed extra repetitions (−1 for infinite number), and #3 tells us about laziness.

```

27495 \cs_new:Npn \__regex_msg_repeated:nnN #1#2#3
27496 {
27497     \str_if_eq:eeF { #1 #2 } { 1 0 }
27498     {
27499         , ~ repeated ~
27500         \int_case:nnF {#2}
27501         {
27502             { -1 } { #1~or-more-times,~\bool_if:NTF #3 { lazy } { greedy } }
27503             { 0 } { #1~times }
27504         }
27505         {
27506             between~#1~and~\int_eval:n {#1+#2}~times,~
27507             \bool_if:NTF #3 { lazy } { greedy }
27508         }
27509     }
27510 }

```

(End definition for __regex_msg_repeated:nnN.)

41.9 Code for tracing

There is a more extensive implementation of tracing in the l3trial package l3trace. Function names are a bit different but could be merged.

__regex_trace_push:nnN Here #1 is the module name (regex) and #2 is typically 1. If the module's current tracing level is less than #2 show nothing, otherwise write #3 to the terminal.

```

\__regex_trace_pop:nnN
\__regex_trace:nnx
27511 \cs_new_protected:Npn \__regex_trace_push:nnN #1#2#3
27512 { \__regex_trace:nnx {#1} {#2} { entering~ \token_to_str:N #3 } }
27513 \cs_new_protected:Npn \__regex_trace_pop:nnN #1#2#3

```

```

27514 { \_regex_trace:nxx {#1} {#2} { leaving~ \token_to_str:N #3 } }
27515 \cs_new_protected:Npn \_regex_trace:nxx #1#2#3
27516 {
27517   \int_compare:nNnF
27518     { \int_use:c { g\_regex_trace_#1_int } } < {#2}
27519     { \iow_term:x { Trace:~#3 } }
27520 }

```

(End definition for _regex_trace_push:nnN, _regex_trace_pop:nnN, and _regex_trace:nxx.)

\g_regex_trace_regex_int No tracing when that is zero.

```

27521 \int_new:N \g_regex_trace_regex_int

```

(End definition for \g_regex_trace_regex_int.)

_regex_trace_states:n This function lists the contents of all states of the NFA, stored in \toks from 0 to \l__regex_max_state_int (excluded).

```

27522 \cs_new_protected:Npn \_regex_trace_states:n #1
27523 {
27524   \int_step_inline:nnn
27525     \l__regex_min_state_int
27526     { \l__regex_max_state_int - 1 }
27527     {
27528       \_regex_trace:nxx { regex } {#1}
27529       { \iow_char:N \\toks ##1 = { \_regex_toks_use:w ##1 } }
27530     }
27531 }

```

(End definition for _regex_trace_states:n.)

```

27532 </package>

```

42 l3box implementation

```

27533 <*package>

```

```

27534 <@@=box>

```

42.1 Support code

_box_dim_eval:w Evaluating a dimension expression expandably. The only difference with \dim_eval:n is the lack of \dim_use:N, to produce an internal dimension rather than expand it into characters.

```

27535 \cs_new_eq:NN \_box_dim_eval:w \tex_dimexpr:D
27536 \cs_new:Npn \_box_dim_eval:n #1
27537 { \_box_dim_eval:w #1 \scan_stop: }

```

(End definition for _box_dim_eval:w and _box_dim_eval:n.)

42.2 Creating and initialising boxes

The following test files are used for this code: *m3box001.lvt*.

\box_new:N Defining a new $\langle box \rangle$ register: remember that box 255 is not generally available.

```
\box_new:c 27538 \cs_new_protected:Npn \box_new:N #1
27539 {
27540     \__kernel_chk_if_free_cs:N #1
27541     \cs:w newbox \cs_end: #1
27542 }
27543 \cs_generate_variant:Nn \box_new:N { c }
```

Clear a $\langle box \rangle$ register.

```
27544 \cs_new_protected:Npn \box_clear:N #1
\box_clear:N 27545 { \box_set_eq:NN #1 \c_empty_box }
27546 \cs_new_protected:Npn \box_gclear:N #1
\box_gclear:N 27547 { \box_gset_eq:NN #1 \c_empty_box }
27548 \cs_generate_variant:Nn \box_clear:N { c }
27549 \cs_generate_variant:Nn \box_gclear:N { c }
```

Clear or new.

```
27550 \cs_new_protected:Npn \box_clear_new:N #1
\box_clear_new:N 27551 { \box_if_exist:NTF #1 { \box_clear:N #1 } { \box_new:N #1 } }
27552 \cs_new_protected:Npn \box_gclear_new:N #1
\box_gclear_new:N 27553 { \box_if_exist:NTF #1 { \box_gclear:N #1 } { \box_new:N #1 } }
27554 \cs_generate_variant:Nn \box_clear_new:N { c }
27555 \cs_generate_variant:Nn \box_gclear_new:N { c }
```

Assigning the contents of a box to be another box.

```
27556 \cs_new_protected:Npn \box_set_eq:NN #1#2
\box_set_eq:NN 27557 { \tex_setbox:D #1 \tex_copy:D #2 }
\box_set_eq:cN 27558 \cs_new_protected:Npn \box_gset_eq:NN #1#2
\box_set_eq:Nc 27559 { \tex_global:D \tex_setbox:D #1 \tex_copy:D #2 }
\box_set_eq:cc 27560 \cs_generate_variant:Nn \box_set_eq:NN { c , Nc , cc }
\box_gset_eq:NN 27561 \cs_generate_variant:Nn \box_gset_eq:NN { c , Nc , cc }
```

Assigning the contents of a box to be another box, then drops the original box.

```
27562 \cs_new_protected:Npn \box_set_eq_drop:NN #1#2
\box_set_eq_drop:NN 27563 { \tex_setbox:D #1 \tex_box:D #2 }
\box_set_eq_drop:cN 27564 \cs_new_protected:Npn \box_gset_eq_drop:NN #1#2
\box_set_eq_drop:Nc 27565 { \tex_global:D \tex_setbox:D #1 \tex_box:D #2 }
\box_set_eq_drop:cc 27566 \cs_generate_variant:Nn \box_set_eq_drop:NN { c , Nc , cc }
\box_gset_eq_drop:NN 27567 \cs_generate_variant:Nn \box_gset_eq_drop:NN { c , Nc , cc }
```

Copies of the cs functions defined in l3basics.

```
27568 \prg_new_eq_conditional:NNn \box_if_exist:N \cs_if_exist:N
\box_if_exist:N 27569 { TF , T , F , p }
\box_if_exist:p:c 27570 \prg_new_eq_conditional:NNn \box_if_exist:c \cs_if_exist:c
\box_if_exist:NTF 27571 { TF , T , F , p }
\box_if_exist:cTF
```

42.3 Measuring and setting box dimensions

Accessing the height, depth, and width of a $\langle box \rangle$ register.

```

27572 \cs_new_eq:NN \box_ht:N \tex_ht:D
\box_ht:N 27573 \cs_new_eq:NN \box_dp:N \tex_dp:D
\box_ht:c 27574 \cs_new_eq:NN \box_wd:N \tex_wd:D
\box_dp:N 27575 \cs_generate_variant:Nn \box_ht:N { c }
\box_dp:c 27576 \cs_generate_variant:Nn \box_dp:N { c }
\box_wd:N 27577 \cs_generate_variant:Nn \box_wd:N { c }
\box_wd:c

```

Setting the size whilst respecting local scope requires copying; the same issue does not come up when working globally. When debugging, the dimension expression #2 is surrounded by parentheses to catch early termination.

```

\box_set_ht:Nn 27578 \cs_new_protected:Npn \box_set_dp:Nn #1#2
\box_set_ht:cn 27579 {
\box_gset_ht:Nn 27580 \tex_setbox:D #1 = \tex_copy:D #1
\box_gset_ht:cn 27581 \box_dp:N #1 \_box_dim_eval:n {#2}
\box_set_dp:Nn 27582 }
\box_set_dp:cn 27583 \cs_generate_variant:Nn \box_set_dp:Nn { c }
\box_gset_dp:Nn 27584 \cs_new_protected:Npn \box_gset_dp:Nn #1#2
\box_gset_dp:cn 27585 { \box_dp:N #1 \_box_dim_eval:n {#2} }
\box_set_wd:Nn 27586 \cs_generate_variant:Nn \box_gset_dp:Nn { c }
\box_set_wd:cn 27587 \cs_new_protected:Npn \box_set_ht:Nn #1#2
\box_gset_wd:Nn 27588 {
\box_gset_wd:cn 27589 \tex_setbox:D #1 = \tex_copy:D #1
27590 \box_ht:N #1 \_box_dim_eval:n {#2}
27591 }
27592 \cs_generate_variant:Nn \box_set_ht:Nn { c }
27593 \cs_new_protected:Npn \box_gset_ht:Nn #1#2
27594 { \box_ht:N #1 \_box_dim_eval:n {#2} }
27595 \cs_generate_variant:Nn \box_gset_ht:Nn { c }
27596 \cs_new_protected:Npn \box_set_wd:Nn #1#2
27597 {
27598 \tex_setbox:D #1 = \tex_copy:D #1
27599 \box_wd:N #1 \_box_dim_eval:n {#2}
27600 }
27601 \cs_generate_variant:Nn \box_set_wd:Nn { c }
27602 \cs_new_protected:Npn \box_gset_wd:Nn #1#2
27603 { \box_wd:N #1 \_box_dim_eval:n {#2} }
27604 \cs_generate_variant:Nn \box_gset_wd:Nn { c }

```

42.4 Using boxes

Using a $\langle box \rangle$. These are just \TeX primitives with meaningful names.

```

27605 \cs_new_eq:NN \box_use_drop:N \tex_box:D
\box_use_drop:N 27606 \cs_new_eq:NN \box_use:N \tex_copy:D
\box_use_drop:c 27607 \cs_generate_variant:Nn \box_use_drop:N { c }
\box_use:N 27608 \cs_generate_variant:Nn \box_use:N { c }
\box_use:c

```

Move box material in different directions. When debugging, the dimension expression #1 is surrounded by parentheses to catch early termination.

```

\box_move_left:nn 27609 \cs_new_protected:Npn \box_move_left:nn #1#2
\box_move_right:nn 27610 { \tex_moveleft:D \_box_dim_eval:n {#1} #2 }
\box_move_up:nn
\box_move_down:nn

```

```

27611 \cs_new_protected:Npn \box_move_right:nn #1#2
27612 { \tex_moveright:D \__box_dim_eval:n {#1} #2 }
27613 \cs_new_protected:Npn \box_move_up:nn #1#2
27614 { \tex_raise:D \__box_dim_eval:n {#1} #2 }
27615 \cs_new_protected:Npn \box_move_down:nn #1#2
27616 { \tex_lower:D \__box_dim_eval:n {#1} #2 }

```

42.5 Box conditionals

The primitives for testing if a $\langle box \rangle$ is empty/void or which type of box it is.

```

\if_hbox:N
\if_vbox:N
\if_box_empty:N

```

```

27617 \cs_new_eq:NN \if_hbox:N \tex_ifhbox:D
27618 \cs_new_eq:NN \if_vbox:N \tex_ifvbox:D
27619 \cs_new_eq:NN \if_box_empty:N \tex_ifvoid:D

```

```

\box_if_horizontal_p:N
\box_if_horizontal_p:c
\box_if_horizontal:N $\overline{TF}$ 
\box_if_horizontal:c $\overline{TF}$ 
\box_if_vertical_p:N
\box_if_vertical_p:c
\box_if_vertical:N $\overline{TF}$ 
\box_if_vertical:c $\overline{TF}$ 

```

```

27620 \prg_new_conditional:Npnn \box_if_horizontal:N #1 { p , T , F , TF }
27621 { \if_hbox:N #1 \prg_return_true: \else: \prg_return_false: \fi: }
27622 \prg_new_conditional:Npnn \box_if_vertical:N #1 { p , T , F , TF }
27623 { \if_vbox:N #1 \prg_return_true: \else: \prg_return_false: \fi: }
27624 \prg_generate_conditional_variant:Nnn \box_if_horizontal:N
27625 { c } { p , T , F , TF }
27626 \prg_generate_conditional_variant:Nnn \box_if_vertical:N
27627 { c } { p , T , F , TF }

```

Testing if a $\langle box \rangle$ is empty/void.

```

\box_if_empty_p:N
\box_if_empty_p:c
\box_if_empty:N $\overline{TF}$ 
\box_if_empty:c $\overline{TF}$ 

```

```

27628 \prg_new_conditional:Npnn \box_if_empty:N #1 { p , T , F , TF }
27629 { \if_box_empty:N #1 \prg_return_true: \else: \prg_return_false: \fi: }
27630 \prg_generate_conditional_variant:Nnn \box_if_empty:N
27631 { c } { p , T , F , TF }

```

(End definition for $\backslash\text{box_new:N}$ and others. These functions are documented on page 242.)

42.6 The last box inserted

```

\box_set_to_last:N
\box_set_to_last:c
\box_gset_to_last:N
\box_gset_to_last:c

```

Set a box to the previous box.

```

27632 \cs_new_protected:Npn \box_set_to_last:N #1
27633 { \tex_setbox:D #1 \tex_lastbox:D }
27634 \cs_new_protected:Npn \box_gset_to_last:N #1
27635 { \tex_global:D \tex_setbox:D #1 \tex_lastbox:D }
27636 \cs_generate_variant:Nn \box_set_to_last:N { c }
27637 \cs_generate_variant:Nn \box_gset_to_last:N { c }

```

(End definition for $\backslash\text{box_set_to_last:N}$ and $\backslash\text{box_gset_to_last:N}$. These functions are documented on page 244.)

42.7 Constant boxes

```

\c_empty_box

```

A box we never use.

```

27638 \box_new:N \c_empty_box

```

(End definition for $\backslash\text{c_empty_box}$. This variable is documented on page 244.)

42.8 Scratch boxes

```
\l_tmpa_box Scratch boxes.
\l_tmpb_box 27639 \box_new:N \l_tmpa_box
\g_tmpa_box 27640 \box_new:N \l_tmpb_box
\g_tmpb_box 27641 \box_new:N \g_tmpa_box
27642 \box_new:N \g_tmpb_box
```

(End definition for `\l_tmpa_box` and others. These variables are documented on page 245.)

42.9 Viewing box contents

TeX's `\showbox` is not really that helpful in many cases, and it is also inconsistent with other L^AT_EX3 show functions as it does not actually shows material in the terminal. So we provide a richer set of functionality.

```
\box_show:N Essentially a wrapper around the internal function, but evaluating the breadth and depth
\box_show:c arguments now outside the group.
\box_show:Nnn 27643 \cs_new_protected:Npn \box_show:N #1
\box_show:cnn 27644 { \box_show:Nnn #1 \c_max_int \c_max_int }
27645 \cs_generate_variant:Nn \box_show:N { c }
27646 \cs_new_protected:Npn \box_show:Nnn #1#2#3
27647 { \__box_show:NNff 1 #1 { \int_eval:n {#2} } { \int_eval:n {#3} } }
27648 \cs_generate_variant:Nn \box_show:Nnn { c }
```

(End definition for `\box_show:N` and `\box_show:Nnn`. These functions are documented on page 245.)

```
\box_log:N Getting TeX to write to the log without interruption the run is done by altering the
\box_log:c interaction mode. For that, the  $\varepsilon$ -TeX extensions are needed.
\box_log:Nnn 27649 \cs_new_protected:Npn \box_log:N #1
\box_log:cnn 27650 { \box_log:Nnn #1 \c_max_int \c_max_int }
\__box_log:nNnn 27651 \cs_generate_variant:Nn \box_log:N { c }
27652 \cs_new_protected:Npn \box_log:Nnn
27653 { \exp_args:No \__box_log:nNnn { \tex_the:D \tex_interactionmode:D } }
27654 \cs_new_protected:Npn \__box_log:nNnn #1#2#3#4
27655 {
27656 \int_set:Nn \tex_interactionmode:D { 0 }
27657 \__box_show:NNff 0 #2 { \int_eval:n {#3} } { \int_eval:n {#4} }
27658 \int_set:Nn \tex_interactionmode:D {#1}
27659 }
27660 \cs_generate_variant:Nn \box_log:Nnn { c }
```

(End definition for `\box_log:N`, `\box_log:Nnn`, and `__box_log:nNnn`. These functions are documented on page 245.)

```
\__box_show:NNnn The internal auxiliary to actually do the output uses a group to deal with breadth and
\__box_show:NNff depth values. The \use:n here gives better output appearance. Setting \tracingonline
and \errorcontextlines is used to control what appears in the terminal.
27661 \cs_new_protected:Npn \__box_show:NNnn #1#2#3#4
27662 {
27663 \box_if_exist:NTF #2
27664 {
27665 \group_begin:
27666 \int_set:Nn \tex_showboxbreadth:D {#3}
```

```

27667         \int_set:Nn \tex_showboxdepth:D    {#4}
27668         \int_set:Nn \tex_tracingonline:D    {#1}
27669         \int_set:Nn \tex_errorcontextlines:D { -1 }
27670         \tex_showbox:D \use:n {#2}
27671     \group_end:
27672 }
27673 {
27674     \__kernel_msg_error:nxx { kernel } { variable-not-defined }
27675     { \token_to_str:N #2 }
27676 }
27677 }
27678 \cs_generate_variant:Nn \__box_show:NNnn { NNff }

```

(End definition for __box_show:NNnn.)

42.10 Horizontal mode boxes

\hbox:n (The test suite for this command, and others in this file, is *m3box002.lvt*.)
Put a horizontal box directly into the input stream.

```

27679 \cs_new_protected:Npn \hbox:n #1
27680 { \tex_hbox:D \scan_stop: { \color_group_begin: #1 \color_group_end: } }

```

(End definition for \hbox:n. This function is documented on page 245.)

```

\hbox_set:Nn
\hbox_set:cn
\hbox_gset:Nn
\hbox_gset:cn
27681 \cs_new_protected:Npn \hbox_set:Nn #1#2
27682 {
27683     \tex_setbox:D #1 \tex_hbox:D
27684     { \color_group_begin: #2 \color_group_end: }
27685 }
27686 \cs_new_protected:Npn \hbox_gset:Nn #1#2
27687 {
27688     \tex_global:D \tex_setbox:D #1 \tex_hbox:D
27689     { \color_group_begin: #2 \color_group_end: }
27690 }
27691 \cs_generate_variant:Nn \hbox_set:Nn { c }
27692 \cs_generate_variant:Nn \hbox_gset:Nn { c }

```

(End definition for \hbox_set:Nn and \hbox_gset:Nn. These functions are documented on page 246.)

\hbox_set_to_wd:Nnn Storing material in a horizontal box with a specified width. Again, put the dimension expression in parentheses when debugging.

```

\hbox_set_to_wd:cn
\hbox_gset_to_wd:Nnn
\hbox_gset_to_wd:cn
27693 \cs_new_protected:Npn \hbox_set_to_wd:Nnn #1#2#3
27694 {
27695     \tex_setbox:D #1 \tex_hbox:D to \__box_dim_eval:n {#2}
27696     { \color_group_begin: #3 \color_group_end: }
27697 }
27698 \cs_new_protected:Npn \hbox_gset_to_wd:Nnn #1#2#3
27699 {
27700     \tex_global:D \tex_setbox:D #1 \tex_hbox:D to \__box_dim_eval:n {#2}
27701     { \color_group_begin: #3 \color_group_end: }
27702 }
27703 \cs_generate_variant:Nn \hbox_set_to_wd:Nnn { c }
27704 \cs_generate_variant:Nn \hbox_gset_to_wd:Nnn { c }

```

(End definition for `\hbox_set_to_wd:Nnn` and `\hbox_gset_to_wd:Nnn`. These functions are documented on page 246.)

<code>\hbox_set:Nw</code>	Storing material in a horizontal box. This type is useful in environment definitions.
<code>\hbox_set:cw</code>	27705 <code>\cs_new_protected:Npn \hbox_set:Nw #1</code>
<code>\hbox_gset:Nw</code>	27706 <code>{</code>
<code>\hbox_gset:cw</code>	27707 <code>\tex_setbox:D #1 \tex_hbox:D</code>
<code>\hbox_set_end:</code>	27708 <code>\c_group_begin_token</code>
<code>\hbox_gset_end:</code>	27709 <code>\color_group_begin:</code>
	27710 <code>}</code>
	27711 <code>\cs_new_protected:Npn \hbox_gset:Nw #1</code>
	27712 <code>{</code>
	27713 <code>\tex_global:D \tex_setbox:D #1 \tex_hbox:D</code>
	27714 <code>\c_group_begin_token</code>
	27715 <code>\color_group_begin:</code>
	27716 <code>}</code>
	27717 <code>\cs_generate_variant:Nn \hbox_set:Nw { c }</code>
	27718 <code>\cs_generate_variant:Nn \hbox_gset:Nw { c }</code>
	27719 <code>\cs_new_protected:Npn \hbox_set_end:</code>
	27720 <code>{</code>
	27721 <code>\color_group_end:</code>
	27722 <code>\c_group_end_token</code>
	27723 <code>}</code>
	27724 <code>\cs_new_eq:NN \hbox_gset_end: \hbox_set_end:</code>

(End definition for `\hbox_set:Nw` and others. These functions are documented on page 246.)

<code>\hbox_set_to_wd:Nnw</code>	Combining the above ideas.
<code>\hbox_set_to_wd:cnw</code>	27725 <code>\cs_new_protected:Npn \hbox_set_to_wd:Nnw #1#2</code>
<code>\hbox_gset_to_wd:Nnw</code>	27726 <code>{</code>
<code>\hbox_gset_to_wd:cnw</code>	27727 <code>\tex_setbox:D #1 \tex_hbox:D to _box_dim_eval:n {#2}</code>
	27728 <code>\c_group_begin_token</code>
	27729 <code>\color_group_begin:</code>
	27730 <code>}</code>
	27731 <code>\cs_new_protected:Npn \hbox_gset_to_wd:Nnw #1#2</code>
	27732 <code>{</code>
	27733 <code>\tex_global:D \tex_setbox:D #1 \tex_hbox:D to _box_dim_eval:n {#2}</code>
	27734 <code>\c_group_begin_token</code>
	27735 <code>\color_group_begin:</code>
	27736 <code>}</code>
	27737 <code>\cs_generate_variant:Nn \hbox_set_to_wd:Nnw { c }</code>
	27738 <code>\cs_generate_variant:Nn \hbox_gset_to_wd:Nnw { c }</code>

(End definition for `\hbox_set_to_wd:Nnw` and `\hbox_gset_to_wd:Nnw`. These functions are documented on page 246.)

<code>\hbox_to_wd:nn</code>	Put a horizontal box directly into the input stream.
<code>\hbox_to_zero:nn</code>	27739 <code>\cs_new_protected:Npn \hbox_to_wd:nn #1#2</code>
	27740 <code>{</code>
	27741 <code>\tex_hbox:D to _box_dim_eval:n {#1}</code>
	27742 <code>{ \color_group_begin: #2 \color_group_end: }</code>
	27743 <code>}</code>
	27744 <code>\cs_new_protected:Npn \hbox_to_zero:nn #1</code>
	27745 <code>{</code>
	27746 <code>\tex_hbox:D to \c_zero_dim</code>

```

27747     { \color_group_begin: #1 \color_group_end: }
27748   }

```

(End definition for `\hbox_to_wd:nn` and `\hbox_to_zero:n`. These functions are documented on page 246.)

`\hbox_overlap_center:n` Put a zero-sized box with the contents pushed against one side (which makes it stick out on the other) directly into the input stream.

```

\hbox_overlap_left:n
\hbox_overlap_right:n
27749 \cs_new_protected:Npn \hbox_overlap_center:n #1
27750 { \hbox_to_zero:n { \tex_hss:D #1 \tex_hss:D } }
27751 \cs_new_protected:Npn \hbox_overlap_left:n #1
27752 { \hbox_to_zero:n { \tex_hss:D #1 } }
27753 \cs_new_protected:Npn \hbox_overlap_right:n #1
27754 { \hbox_to_zero:n { #1 \tex_hss:D } }

```

(End definition for `\hbox_overlap_center:n`, `\hbox_overlap_left:n`, and `\hbox_overlap_right:n`. These functions are documented on page 246.)

```

\hbox_unpack:N Unpacking a box and if requested also clear it.
\hbox_unpack:c
\hbox_unpack_drop:N
\hbox_unpack_drop:c
27755 \cs_new_eq:NN \hbox_unpack:N \tex_unhcopy:D
27756 \cs_new_eq:NN \hbox_unpack_drop:N \tex_unhbox:D
27757 \cs_generate_variant:Nn \hbox_unpack:N { c }
27758 \cs_generate_variant:Nn \hbox_unpack_drop:N { c }

```

(End definition for `\hbox_unpack:N` and `\hbox_unpack_drop:N`. These functions are documented on page 247.)

42.11 Vertical mode boxes

TeX ends these boxes directly with the internal `end_graf` routine. This means that there is no `\par` at the end of vertical boxes unless we insert one. Thus all vertical boxes include a `\par` just before closing the color group.

`\vbox:n` The following test files are used for this code: `m3box003.lvt`.

The following test files are used for this code: `m3box003.lvt`.

```

\vbox_top:n Put a vertical box directly into the input stream.
27759 \cs_new_protected:Npn \vbox:n #1
27760 { \tex_vbox:D { \color_group_begin: #1 \par \color_group_end: } }
27761 \cs_new_protected:Npn \vbox_top:n #1
27762 { \tex_vtop:D { \color_group_begin: #1 \par \color_group_end: } }

```

(End definition for `\vbox:n` and `\vbox_top:n`. These functions are documented on page 247.)

```

\vbox_to_ht:nn Put a vertical box directly into the input stream.
\vbox_to_zero:n
\vbox_to_ht:nn
\vbox_to_zero:n
27763 \cs_new_protected:Npn \vbox_to_ht:nn #1#2
27764 {
27765   \tex_vbox:D to \__box_dim_eval:n {#1}
27766   { \color_group_begin: #2 \par \color_group_end: }
27767 }
27768 \cs_new_protected:Npn \vbox_to_zero:n #1
27769 {
27770   \tex_vbox:D to \c_zero_dim
27771   { \color_group_begin: #1 \par \color_group_end: }
27772 }

```

(End definition for `\vbox_to_ht:nn` and others. These functions are documented on page 247.)

`\vbox_set:Nn` Storing material in a vertical box with a natural height.
`\vbox_set:cn` 27773 `\cs_new_protected:Npn \vbox_set:Nn #1#2`
`\vbox_gset:Nn` 27774 `{`
`\vbox_gset:cn` 27775 `\tex_setbox:D #1 \tex_vbox:D`
27776 `{ \color_group_begin: #2 \par \color_group_end: }`
27777 `}`
27778 `\cs_new_protected:Npn \vbox_gset:Nn #1#2`
27779 `{`
27780 `\tex_global:D \tex_setbox:D #1 \tex_vbox:D`
27781 `{ \color_group_begin: #2 \par \color_group_end: }`
27782 `}`
27783 `\cs_generate_variant:Nn \vbox_set:Nn { c }`
27784 `\cs_generate_variant:Nn \vbox_gset:Nn { c }`

(End definition for `\vbox_set:Nn` and `\vbox_gset:Nn`. These functions are documented on page 247.)

`\vbox_set_top:Nn` Storing material in a vertical box with a natural height and reference point at the baseline
`\vbox_set_top:cn` of the first object in the box.
`\vbox_gset_top:Nn` 27785 `\cs_new_protected:Npn \vbox_set_top:Nn #1#2`
`\vbox_gset_top:cn` 27786 `{`
27787 `\tex_setbox:D #1 \tex_vtop:D`
27788 `{ \color_group_begin: #2 \par \color_group_end: }`
27789 `}`
27790 `\cs_new_protected:Npn \vbox_gset_top:Nn #1#2`
27791 `{`
27792 `\tex_global:D \tex_setbox:D #1 \tex_vtop:D`
27793 `{ \color_group_begin: #2 \par \color_group_end: }`
27794 `}`
27795 `\cs_generate_variant:Nn \vbox_set_top:Nn { c }`
27796 `\cs_generate_variant:Nn \vbox_gset_top:Nn { c }`

(End definition for `\vbox_set_top:Nn` and `\vbox_gset_top:Nn`. These functions are documented on page 247.)

`\vbox_set_to_ht:Nnn` Storing material in a vertical box with a specified height.
`\vbox_set_to_ht:cn` 27797 `\cs_new_protected:Npn \vbox_set_to_ht:Nnn #1#2#3`
`\vbox_gset_to_ht:Nnn` 27798 `{`
`\vbox_gset_to_ht:cn` 27799 `\tex_setbox:D #1 \tex_vbox:D to _box_dim_eval:n {#2}`
27800 `{ \color_group_begin: #3 \par \color_group_end: }`
27801 `}`
27802 `\cs_new_protected:Npn \vbox_gset_to_ht:Nnn #1#2#3`
27803 `{`
27804 `\tex_global:D \tex_setbox:D #1 \tex_vbox:D to _box_dim_eval:n {#2}`
27805 `{ \color_group_begin: #3 \par \color_group_end: }`
27806 `}`
27807 `\cs_generate_variant:Nn \vbox_set_to_ht:Nnn { c }`
27808 `\cs_generate_variant:Nn \vbox_gset_to_ht:Nnn { c }`

(End definition for `\vbox_set_to_ht:Nnn` and `\vbox_gset_to_ht:Nnn`. These functions are documented on page 248.)

`\vbox_set:Nw` Storing material in a vertical box. This type is useful in environment definitions.

```

\vbox_set:cnw 27809 \cs_new_protected:Npn \vbox_set:Nw #1
\vbox_gset:Nw 27810 {
\vbox_gset:cnw 27811 \tex_setbox:D #1 \tex_vbox:D
\vbox_set_end: 27812 \c_group_begin_token
\vbox_gset_end: 27813 \color_group_begin:
27814 }
27815 \cs_new_protected:Npn \vbox_gset:Nw #1
27816 {
27817 \tex_global:D \tex_setbox:D #1 \tex_vbox:D
27818 \c_group_begin_token
27819 \color_group_begin:
27820 }
27821 \cs_generate_variant:Nn \vbox_set:Nw { c }
27822 \cs_generate_variant:Nn \vbox_gset:Nw { c }
27823 \cs_new_protected:Npn \vbox_set_end:
27824 {
27825 \par
27826 \color_group_end:
27827 \c_group_end_token
27828 }
27829 \cs_new_eq:NN \vbox_gset_end: \vbox_set_end:

```

(End definition for `\vbox_set:Nw` and others. These functions are documented on page 248.)

`\vbox_set_to_ht:Nnw` A combination of the above ideas.

```

\vbox_set_to_ht:cnw 27830 \cs_new_protected:Npn \vbox_set_to_ht:Nnw #1#2
\vbox_gset_to_ht:Nnw 27831 {
\vbox_gset_to_ht:cnw 27832 \tex_setbox:D #1 \tex_vbox:D to \_box_dim_eval:n {#2}
27833 \c_group_begin_token
27834 \color_group_begin:
27835 }
27836 \cs_new_protected:Npn \vbox_gset_to_ht:Nnw #1#2
27837 {
27838 \tex_global:D \tex_setbox:D #1 \tex_vbox:D to \_box_dim_eval:n {#2}
27839 \c_group_begin_token
27840 \color_group_begin:
27841 }
27842 \cs_generate_variant:Nn \vbox_set_to_ht:Nnw { c }
27843 \cs_generate_variant:Nn \vbox_gset_to_ht:Nnw { c }

```

(End definition for `\vbox_set_to_ht:Nnw` and `\vbox_gset_to_ht:Nnw`. These functions are documented on page 248.)

`\vbox_unpack:N` Unpacking a box and if requested also clear it.

```

\vbox_unpack:c 27844 \cs_new_eq:NN \vbox_unpack:N \tex_unvcopy:D
\vbox_unpack_drop:N 27845 \cs_new_eq:NN \vbox_unpack_drop:N \tex_unvbox:D
\vbox_unpack_drop:c 27846 \cs_generate_variant:Nn \vbox_unpack:N { c }
27847 \cs_generate_variant:Nn \vbox_unpack_drop:N { c }

```

(End definition for `\vbox_unpack:N` and `\vbox_unpack_drop:N`. These functions are documented on page 248.)

```

\ vbox_set_split_to_ht:NNn Splitting a vertical box in two.
\ vbox_set_split_to_ht:cNn
\ vbox_set_split_to_ht:Ncn
\ vbox_set_split_to_ht:ccn
\ vbox_gset_split_to_ht:NNn
\ vbox_gset_split_to_ht:cNn
\ vbox_gset_split_to_ht:Ncn
\ vbox_gset_split_to_ht:ccn
27848 \cs_new_protected:Npn \vbox_set_split_to_ht:NNn #1#2#3
27849 { \tex_setbox:D #1 \tex_vsplit:D #2 to \_box_dim_eval:n {#3} }
27850 \cs_generate_variant:Nn \vbox_set_split_to_ht:NNn { c , Nc , cc }
27851 \cs_new_protected:Npn \vbox_gset_split_to_ht:NNn #1#2#3
27852 {
27853   \tex_global:D \tex_setbox:D #1
27854   \tex_vsplit:D #2 to \_box_dim_eval:n {#3}
27855 }
27856 \cs_generate_variant:Nn \vbox_gset_split_to_ht:NNn { c , Nc , cc }

```

(End definition for `\vbox_set_split_to_ht:NNn` and `\vbox_gset_split_to_ht:NNn`. These functions are documented on page 248.)

42.12 Affine transformations

`\l__box_angle_fp` When rotating boxes, the angle itself may be needed by the engine-dependent code. This is done using the `fp` module so that the value is tidied up properly.

```
27857 \fp_new:N \l__box_angle_fp
```

(End definition for `\l__box_angle_fp`.)

`\l__box_cos_fp` These are used to hold the calculated sine and cosine values while carrying out a rotation.

```
\l__box_sin_fp 27858 \fp_new:N \l__box_cos_fp
```

```
27859 \fp_new:N \l__box_sin_fp
```

(End definition for `\l__box_cos_fp` and `\l__box_sin_fp`.)

`\l__box_top_dim` These are the positions of the four edges of a box before manipulation.

```

\l__box_bottom_dim 27860 \dim_new:N \l__box_top_dim
\l__box_left_dim 27861 \dim_new:N \l__box_bottom_dim
\l__box_right_dim 27862 \dim_new:N \l__box_left_dim
27863 \dim_new:N \l__box_right_dim

```

(End definition for `\l__box_top_dim` and others.)

`\l__box_top_new_dim` These are the positions of the four edges of a box after manipulation.

```

\l__box_bottom_new_dim 27864 \dim_new:N \l__box_top_new_dim
\l__box_left_new_dim 27865 \dim_new:N \l__box_bottom_new_dim
\l__box_right_new_dim 27866 \dim_new:N \l__box_left_new_dim
27867 \dim_new:N \l__box_right_new_dim

```

(End definition for `\l__box_top_new_dim` and others.)

`\l__box_internal_box` Scratch space, but also needed by some parts of the driver.

```
27868 \box_new:N \l__box_internal_box
```

(End definition for `\l__box_internal_box`.)

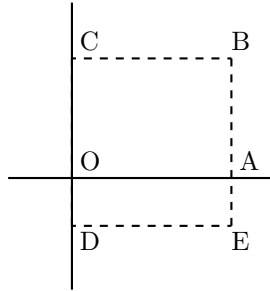


Figure 1: Co-ordinates of a box prior to rotation.

\box_rotate:Nn Rotation of a box starts with working out the relevant sine and cosine. The actual
\box_rotate:cn rotation is in an auxiliary to keep the flow slightly clearer

\box_grotate:Nn 27869 \cs_new_protected:Npn \box_rotate:Nn #1#2

\box_grotate:cn 27870 { __box_rotate:NnN #1 {#2} \hbox_set:Nn }

__box_rotate:NnN 27871 \cs_generate_variant:Nn \box_rotate:Nn { c }

__box_rotate:N 27872 \cs_new_protected:Npn \box_grotate:Nn #1#2

__box_rotate_xdir:nnN 27873 { __box_rotate:NnN #1 {#2} \hbox_gset:Nn }

__box_rotate_ydir:nnN 27874 \cs_generate_variant:Nn \box_grotate:Nn { c }

__box_rotate_quadrant_one: 27875 \cs_new_protected:Npn __box_rotate:NnN #1#2#3

__box_rotate_quadrant_two: 27876 {

__box_rotate_quadrant_three: 27877 #3 #1

__box_rotate_quadrant_four: 27878 {

27879 \fp_set:Nn \l__box_angle_fp {#2}

27880 \fp_set:Nn \l__box_sin_fp { sind (\l__box_angle_fp) }

27881 \fp_set:Nn \l__box_cos_fp { cosd (\l__box_angle_fp) }

27882 __box_rotate:N #1

27883 }

27884 }

The edges of the box are then recorded: the left edge is always at zero. Rotation of the four edges then takes place: this is most efficiently done on a quadrant by quadrant basis.

27885 \cs_new_protected:Npn __box_rotate:N #1

27886 {

27887 \dim_set:Nn \l__box_top_dim { \box_ht:N #1 }

27888 \dim_set:Nn \l__box_bottom_dim { -\box_dp:N #1 }

27889 \dim_set:Nn \l__box_right_dim { \box_wd:N #1 }

27890 \dim_zero:N \l__box_left_dim

The next step is to work out the x and y coordinates of vertices of the rotated box in relation to its original coordinates. The box can be visualized with vertices B , C , D and E is illustrated (Figure 1). The vertex O is the reference point on the baseline, and in this implementation is also the centre of rotation. The formulae are, for a point P and angle α :

$$\begin{aligned} P'_x &= P_x - O_x \\ P'_y &= P_y - O_y \\ P''_x &= (P'_x \cos(\alpha)) - (P'_y \sin(\alpha)) \\ P''_y &= (P'_x \sin(\alpha)) + (P'_y \cos(\alpha)) \\ P'''_x &= P''_x + O_x + L_x \\ P'''_y &= P''_y + O_y \end{aligned}$$

The “extra” horizontal translation L_x at the end is calculated so that the leftmost point of the resulting box has x -coordinate 0. This is desirable as \TeX boxes must have the reference point at the left edge of the box. (As O is always $(0,0)$, this part of the calculation is omitted here.)

```

27891 \fp_compare:nNnTF \l__box_sin_fp > \c_zero_fp
27892 {
27893   \fp_compare:nNnTF \l__box_cos_fp > \c_zero_fp
27894   { \__box_rotate_quadrant_one: }
27895   { \__box_rotate_quadrant_two: }
27896 }
27897 {
27898   \fp_compare:nNnTF \l__box_cos_fp < \c_zero_fp
27899   { \__box_rotate_quadrant_three: }
27900   { \__box_rotate_quadrant_four: }
27901 }

```

The position of the box edges are now known, but the box at this stage be misplaced relative to the current \TeX reference point. So the content of the box is moved such that the reference point of the rotated box is in the same place as the original.

```

27902 \hbox_set:Nn \l__box_internal_box { \box_use:N #1 }
27903 \hbox_set:Nn \l__box_internal_box
27904 {
27905   \tex_kern:D -\l__box_left_new_dim
27906   \hbox:n
27907   {
27908     \__box_backend_rotate:Nn
27909     \l__box_internal_box
27910     \l__box_angle_fp
27911   }
27912 }

```

Tidy up the size of the box so that the material is actually inside the bounding box. The result can then be used to reset the original box.

```

27913 \box_set_ht:Nn \l__box_internal_box { \l__box_top_new_dim }
27914 \box_set_dp:Nn \l__box_internal_box { -\l__box_bottom_new_dim }
27915 \box_set_wd:Nn \l__box_internal_box
27916 { \l__box_right_new_dim - \l__box_left_new_dim }
27917 \box_use_drop:N \l__box_internal_box
27918 }

```

These functions take a general point $(\#1,\#2)$ and rotate its location about the origin, using the previously-set sine and cosine values. Each function gives only one component of the location of the updated point. This is because for rotation of a box each step needs only one value, and so performance is gained by avoiding working out both x' and y' at the same time. Contrast this with the equivalent function in the `l3coffins` module, where both parts are needed.

```

27919 \cs_new_protected:Npn \__box_rotate_xdir:nnN #1#2#3
27920 {
27921   \dim_set:Nn #3
27922   {
27923     \fp_to_dim:n
27924     {
27925       \l__box_cos_fp * \dim_to_fp:n {#1}
27926       - \l__box_sin_fp * \dim_to_fp:n {#2}

```

```

27927     }
27928   }
27929 }
27930 \cs_new_protected:Npn \__box_rotate_ydir:nnN #1#2#3
27931 {
27932   \dim_set:Nn #3
27933   {
27934     \fp_to_dim:n
27935     {
27936       \l__box_sin_fp * \dim_to_fp:n {#1}
27937       + \l__box_cos_fp * \dim_to_fp:n {#2}
27938     }
27939   }
27940 }

```

Rotation of the edges is done using a different formula for each quadrant. In every case, the top and bottom edges only need the resulting y -values, whereas the left and right edges need the x -values. Each case is a question of picking out which corner ends up at with the maximum top, bottom, left and right value. Doing this by hand means a lot less calculating and avoids lots of comparisons.

```

27941 \cs_new_protected:Npn \__box_rotate_quadrant_one:
27942 {
27943   \__box_rotate_ydir:nnN \l__box_right_dim \l__box_top_dim
27944   \l__box_top_new_dim
27945   \__box_rotate_ydir:nnN \l__box_left_dim \l__box_bottom_dim
27946   \l__box_bottom_new_dim
27947   \__box_rotate_xdir:nnN \l__box_left_dim \l__box_top_dim
27948   \l__box_left_new_dim
27949   \__box_rotate_xdir:nnN \l__box_right_dim \l__box_bottom_dim
27950   \l__box_right_new_dim
27951 }
27952 \cs_new_protected:Npn \__box_rotate_quadrant_two:
27953 {
27954   \__box_rotate_ydir:nnN \l__box_right_dim \l__box_bottom_dim
27955   \l__box_top_new_dim
27956   \__box_rotate_ydir:nnN \l__box_left_dim \l__box_top_dim
27957   \l__box_bottom_new_dim
27958   \__box_rotate_xdir:nnN \l__box_right_dim \l__box_top_dim
27959   \l__box_left_new_dim
27960   \__box_rotate_xdir:nnN \l__box_left_dim \l__box_bottom_dim
27961   \l__box_right_new_dim
27962 }
27963 \cs_new_protected:Npn \__box_rotate_quadrant_three:
27964 {
27965   \__box_rotate_ydir:nnN \l__box_left_dim \l__box_bottom_dim
27966   \l__box_top_new_dim
27967   \__box_rotate_ydir:nnN \l__box_right_dim \l__box_top_dim
27968   \l__box_bottom_new_dim
27969   \__box_rotate_xdir:nnN \l__box_right_dim \l__box_bottom_dim
27970   \l__box_left_new_dim
27971   \__box_rotate_xdir:nnN \l__box_left_dim \l__box_top_dim
27972   \l__box_right_new_dim
27973 }
27974 \cs_new_protected:Npn \__box_rotate_quadrant_four:

```

```

27975 {
27976   \__box_rotate_ydir:nnN \l__box_left_dim \l__box_top_dim
27977   \l__box_top_new_dim
27978   \__box_rotate_ydir:nnN \l__box_right_dim \l__box_bottom_dim
27979   \l__box_bottom_new_dim
27980   \__box_rotate_xdir:nnN \l__box_left_dim \l__box_bottom_dim
27981   \l__box_left_new_dim
27982   \__box_rotate_xdir:nnN \l__box_right_dim \l__box_top_dim
27983   \l__box_right_new_dim
27984 }

```

(End definition for `\box_rotate:Nn` and others. These functions are documented on page 252.)

`\l__box_scale_x_fp` Scaling is potentially-different in the two axes.

```

\l__box_scale_y_fp
27985 \fp_new:N \l__box_scale_x_fp
27986 \fp_new:N \l__box_scale_y_fp

```

(End definition for `\l__box_scale_x_fp` and `\l__box_scale_y_fp`.)

`\box_resize_to_wd_and_ht_plus_dp:Nnn` Resizing a box starts by working out the various dimensions of the existing box.

```

\box_resize_to_wd_and_ht_plus_dp:cnm
27987 \cs_new_protected:Npn \box_resize_to_wd_and_ht_plus_dp:Nnn #1#2#3
\box_gresize_to_wd_and_ht_plus_dp:Nnn
27988 {
\box_gresize_to_wd_and_ht_plus_dp:cnm
27989   \__box_resize_to_wd_and_ht_plus_dp:NnnN #1 {#2} {#3}
\__box_resize_to_wd_and_ht_plus_dp:NnnN
27990   \hbox_set:Nn
\__box_resize_set_corners:N
27991 }
\__box_resize:N
27992 \cs_generate_variant:Nn \box_resize_to_wd_and_ht_plus_dp:Nnn { c }
\__box_resize:NNN
27993 \cs_new_protected:Npn \box_gresize_to_wd_and_ht_plus_dp:Nnn #1#2#3
27994 {
27995   \__box_resize_to_wd_and_ht_plus_dp:NnnN #1 {#2} {#3}
27996   \hbox_gset:Nn
27997 }
27998 \cs_generate_variant:Nn \box_gresize_to_wd_and_ht_plus_dp:Nnn { c }
27999 \cs_new_protected:Npn \__box_resize_to_wd_and_ht_plus_dp:NnnN #1#2#3#4
28000 {
28001   #4 #1
28002   {
28003     \__box_resize_set_corners:N #1

```

The x -scaling and resulting box size is easy enough to work out: the dimension is that given as #2, and the scale is simply the new width divided by the old one.

```

28004   \fp_set:Nn \l__box_scale_x_fp
28005   { \dim_to_fp:n {#2} / \dim_to_fp:n { \l__box_right_dim } }

```

The y -scaling needs both the height and the depth of the current box.

```

28006   \fp_set:Nn \l__box_scale_y_fp
28007   {
28008     \dim_to_fp:n {#3}
28009     / \dim_to_fp:n { \l__box_top_dim - \l__box_bottom_dim }
28010   }

```

Hand off to the auxiliary which does the rest of the work.

```

28011   \__box_resize:N #1
28012 }
28013 }
28014 \cs_new_protected:Npn \__box_resize_set_corners:N #1
28015 {

```

```

28016 \dim_set:Nn \l__box_top_dim { \box_ht:N #1 }
28017 \dim_set:Nn \l__box_bottom_dim { -\box_dp:N #1 }
28018 \dim_set:Nn \l__box_right_dim { \box_wd:N #1 }
28019 \dim_zero:N \l__box_left_dim
28020 }

```

With at least one real scaling to do, the next phase is to find the new edge co-ordinates. In the x direction this is relatively easy: just scale the right edge. In the y direction, both dimensions have to be scaled, and this again needs the absolute scale value. Once that is all done, the common resize/rescale code can be employed.

```

28021 \cs_new_protected:Npn \__box_resize:N #1
28022 {
28023   \__box_resize:NNN \l__box_right_new_dim
28024   \l__box_scale_x_fp \l__box_right_dim
28025   \__box_resize:NNN \l__box_bottom_new_dim
28026   \l__box_scale_y_fp \l__box_bottom_dim
28027   \__box_resize:NNN \l__box_top_new_dim
28028   \l__box_scale_y_fp \l__box_top_dim
28029   \__box_resize_common:N #1
28030 }
28031 \cs_new_protected:Npn \__box_resize:NNN #1#2#3
28032 {
28033   \dim_set:Nn #1
28034   { \fp_to_dim:n { \fp_abs:n { #2 } * \dim_to_fp:n { #3 } } }
28035 }

```

(End definition for `\box_resize_to_wd_and_ht_plus_dp:Nnn` and others. These functions are documented on page 251.)

```

\box_resize_to_ht:Nn
\box_resize_to_ht:cn
\box_gresize_to_ht:Nn
\box_gresize_to_ht:cn
\__box_resize_to_ht:NnN
\box_resize_to_ht_plus_dp:Nn
\box_resize_to_ht_plus_dp:cn
\box_gresize_to_ht_plus_dp:Nn
\box_gresize_to_ht_plus_dp:cn
\__box_resize_to_ht_plus_dp:NnN
\box_resize_to_wd:Nn
\box_resize_to_wd:cn
\box_gresize_to_wd:Nn
\box_gresize_to_wd:cn
\__box_resize_to_wd:NnN
\box_resize_to_wd_and_ht:Nnn
\box_resize_to_wd_and_ht:cn
\box_gresize_to_wd_and_ht:Nnn
\box_gresize_to_wd_and_ht:cn
\__box_resize_to_wd_and_ht:NnnN

```

Scaling to a (total) height or to a width is a simplified version of the main resizing operation, with the scale simply copied between the two parts. The internal auxiliary is called using the scaling value twice, as the sign for both parts is needed (as this allows the same internal code to be used as for the general case).

```

28036 \cs_new_protected:Npn \box_resize_to_ht:Nn #1#2
28037 { \__box_resize_to_ht:NnN #1 {#2} \hbox_set:Nn }
28038 \cs_generate_variant:Nn \box_resize_to_ht:Nn { c }
28039 \cs_new_protected:Npn \box_gresize_to_ht:Nn #1#2
28040 { \__box_resize_to_ht:NnN #1 {#2} \hbox_gset:Nn }
28041 \cs_generate_variant:Nn \box_gresize_to_ht:Nn { c }
28042 \cs_new_protected:Npn \__box_resize_to_ht:NnN #1#2#3
28043 {
28044   #3 #1
28045   {
28046     \__box_resize_set_corners:N #1
28047     \fp_set:Nn \l__box_scale_y_fp
28048     {
28049       \dim_to_fp:n {#2}
28050       / \dim_to_fp:n { \l__box_top_dim }
28051     }
28052     \fp_set_eq:NN \l__box_scale_x_fp \l__box_scale_y_fp
28053     \__box_resize:N #1
28054   }
28055 }
28056 \cs_new_protected:Npn \box_resize_to_ht_plus_dp:Nn #1#2

```

```

28057 { \_box_resize_to_ht_plus_dp:NnN #1 {#2} \hbox_set:Nn }
28058 \cs_generate_variant:Nn \box_resize_to_ht_plus_dp:Nn { c }
28059 \cs_new_protected:Npn \box_gresize_to_ht_plus_dp:Nn #1#2
28060 { \_box_resize_to_ht_plus_dp:NnN #1 {#2} \hbox_gset:Nn }
28061 \cs_generate_variant:Nn \box_gresize_to_ht_plus_dp:Nn { c }
28062 \cs_new_protected:Npn \_box_resize_to_ht_plus_dp:NnN #1#2#3
28063 {
28064   \hbox_set:Nn #1
28065   {
28066     \_box_resize_set_corners:N #1
28067     \fp_set:Nn \l__box_scale_y_fp
28068     {
28069       \dim_to_fp:n {#2}
28070       / \dim_to_fp:n { \l__box_top_dim - \l__box_bottom_dim }
28071     }
28072     \fp_set_eq:NN \l__box_scale_x_fp \l__box_scale_y_fp
28073     \_box_resize:N #1
28074   }
28075 }
28076 \cs_new_protected:Npn \box_resize_to_wd:Nn #1#2
28077 { \_box_resize_to_wd:NnN #1 {#2} \hbox_set:Nn }
28078 \cs_generate_variant:Nn \box_resize_to_wd:Nn { c }
28079 \cs_new_protected:Npn \box_gresize_to_wd:Nn #1#2
28080 { \_box_resize_to_wd:NnN #1 {#2} \hbox_gset:Nn }
28081 \cs_generate_variant:Nn \box_gresize_to_wd:Nn { c }
28082 \cs_new_protected:Npn \_box_resize_to_wd:NnN #1#2#3
28083 {
28084   #3 #1
28085   {
28086     \_box_resize_set_corners:N #1
28087     \fp_set:Nn \l__box_scale_x_fp
28088     { \dim_to_fp:n {#2} / \dim_to_fp:n { \l__box_right_dim } }
28089     \fp_set_eq:NN \l__box_scale_y_fp \l__box_scale_x_fp
28090     \_box_resize:N #1
28091   }
28092 }
28093 \cs_new_protected:Npn \box_resize_to_wd_and_ht:Nnn #1#2#3
28094 { \_box_resize_to_wd_and_ht:NnnN #1 {#2} {#3} \hbox_set:Nn }
28095 \cs_generate_variant:Nn \box_resize_to_wd_and_ht:Nnn { c }
28096 \cs_new_protected:Npn \box_gresize_to_wd_and_ht:Nnn #1#2#3
28097 { \_box_resize_to_wd_and_ht:NnnN #1 {#2} {#3} \hbox_gset:Nn }
28098 \cs_generate_variant:Nn \box_gresize_to_wd_and_ht:Nnn { c }
28099 \cs_new_protected:Npn \_box_resize_to_wd_and_ht:NnnN #1#2#3#4
28100 {
28101   #4 #1
28102   {
28103     \_box_resize_set_corners:N #1
28104     \fp_set:Nn \l__box_scale_x_fp
28105     { \dim_to_fp:n {#2} / \dim_to_fp:n { \l__box_right_dim } }
28106     \fp_set:Nn \l__box_scale_y_fp
28107     {
28108       \dim_to_fp:n {#3}
28109       / \dim_to_fp:n { \l__box_top_dim }
28110     }

```



```

28111         \_box_resize:N #1
28112     }
28113 }

```

(End definition for `\box_resize_to_ht:Nn` and others. These functions are documented on page 250.)

\box_scale:Nnn When scaling a box, setting the scaling itself is easy enough. The new dimensions are also relatively easy to find, allowing only for the need to keep them positive in all cases. **\box_scale:cnn** Once that is done then after a check for the trivial scaling a hand-off can be made to the common code. **\box_gscale:Nnn** The code here is split into two as this allows sharing with the auto-resizing functions. **\box_gscale:cnn**

```

\__box_scale:NnnN
\__box_scale:N
28114 \cs_new_protected:Npn \box_scale:Nnn #1#2#3
28115 { \__box_scale:NnnN #1 {#2} {#3} \hbox_set:Nn }
28116 \cs_generate_variant:Nn \box_scale:Nnn { c }
28117 \cs_new_protected:Npn \box_gscale:Nnn #1#2#3
28118 { \__box_scale:NnnN #1 {#2} {#3} \hbox_gset:Nn }
28119 \cs_generate_variant:Nn \box_gscale:Nnn { c }
28120 \cs_new_protected:Npn \__box_scale:NnnN #1#2#3#4
28121 {
28122     #4 #1
28123     {
28124         \fp_set:Nn \l__box_scale_x_fp {#2}
28125         \fp_set:Nn \l__box_scale_y_fp {#3}
28126         \__box_scale:N #1
28127     }
28128 }
28129 \cs_new_protected:Npn \__box_scale:N #1
28130 {
28131     \dim_set:Nn \l__box_top_dim { \box_ht:N #1 }
28132     \dim_set:Nn \l__box_bottom_dim { -\box_dp:N #1 }
28133     \dim_set:Nn \l__box_right_dim { \box_wd:N #1 }
28134     \dim_zero:N \l__box_left_dim
28135     \dim_set:Nn \l__box_top_new_dim
28136     { \fp_abs:n { \l__box_scale_y_fp } \l__box_top_dim }
28137     \dim_set:Nn \l__box_bottom_new_dim
28138     { \fp_abs:n { \l__box_scale_y_fp } \l__box_bottom_dim }
28139     \dim_set:Nn \l__box_right_new_dim
28140     { \fp_abs:n { \l__box_scale_x_fp } \l__box_right_dim }
28141     \__box_resize_common:N #1
28142 }

```

(End definition for `\box_scale:Nnn` and others. These functions are documented on page 252.)

\box_autosize_to_wd_and_ht:Nnn Although autosizing a box uses dimensions, it has more in common in implementation with scaling. As such, most of the real work here is done elsewhere. **\box_autosize_to_wd_and_ht:cnn**

```

\box_gautosize_to_wd_and_ht:Nnn
\box_gautosize_to_wd_and_ht:cnn
28143 \cs_new_protected:Npn \box_autosize_to_wd_and_ht:Nnn #1#2#3
28144 { \__box_autosize:NnnN #1 {#2} {#3} { \box_ht:N #1 } \hbox_set:Nn }
28145 \cs_generate_variant:Nn \box_autosize_to_wd_and_ht:Nnn { c }
28146 \cs_new_protected:Npn \box_gautosize_to_wd_and_ht:Nnn #1#2#3
28147 { \__box_autosize:NnnN #1 {#2} {#3} { \box_ht:N #1 } \hbox_gset:Nn }
28148 \cs_generate_variant:Nn \box_gautosize_to_wd_and_ht:Nnn { c }
28149 \cs_new_protected:Npn \box_autosize_to_wd_and_ht_plus_dp:Nnn #1#2#3
28150 {
28151     \__box_autosize:NnnN #1 {#2} {#3} { \box_ht:N #1 + \box_dp:N #1 }

```

```

28152     \hbox_set:Nn
28153   }
28154   \cs_generate_variant:Nn \box_autosize_to_wd_and_ht_plus_dp:Nnn { c }
28155   \cs_new_protected:Npn \box_gautosize_to_wd_and_ht_plus_dp:Nnn #1#2#3
28156   {
28157     \__box_autosize:NnnN #1 {#2} {#3} { \box_ht:N #1 + \box_dp:N #1 }
28158     \hbox_gset:Nn
28159   }
28160   \cs_generate_variant:Nn \box_gautosize_to_wd_and_ht_plus_dp:Nnn { c }
28161   \cs_new_protected:Npn \__box_autosize:NnnN #1#2#3#4#5
28162   {
28163     #5 #1
28164     {
28165       \fp_set:Nn \l__box_scale_x_fp { ( #2 ) / \box_wd:N #1 }
28166       \fp_set:Nn \l__box_scale_y_fp { ( #3 ) / ( #4 ) }
28167       \fp_compare:nNnTF \l__box_scale_x_fp > \l__box_scale_y_fp
28168       { \fp_set_eq:NN \l__box_scale_x_fp \l__box_scale_y_fp }
28169       { \fp_set_eq:NN \l__box_scale_y_fp \l__box_scale_x_fp }
28170       \__box_scale:N #1
28171     }
28172   }

```

(End definition for `\box_autosize_to_wd_and_ht:Nnn` and others. These functions are documented on page 250.)

`__box_resize_common:N` The main resize function places its input into a box which start off with zero width, and includes the handles for engine rescaling.

```

28173   \cs_new_protected:Npn \__box_resize_common:N #1
28174   {
28175     \hbox_set:Nn \l__box_internal_box
28176     {
28177       \__box_backend_scale:Nnn
28178       #1
28179       \l__box_scale_x_fp
28180       \l__box_scale_y_fp
28181     }

```

The new height and depth can be applied directly.

```

28182     \fp_compare:nNnTF \l__box_scale_y_fp > \c_zero_fp
28183     {
28184       \box_set_ht:Nn \l__box_internal_box { \l__box_top_new_dim }
28185       \box_set_dp:Nn \l__box_internal_box { -\l__box_bottom_new_dim }
28186     }
28187     {
28188       \box_set_dp:Nn \l__box_internal_box { \l__box_top_new_dim }
28189       \box_set_ht:Nn \l__box_internal_box { -\l__box_bottom_new_dim }
28190     }

```

Things are not quite as obvious for the width, as the reference point needs to remain unchanged. For positive scaling factors resizing the box is all that is needed. However, for case of a negative scaling the material must be shifted such that the reference point ends up in the right place.

```

28191     \fp_compare:nNnTF \l__box_scale_x_fp < \c_zero_fp
28192     {
28193       \hbox_to_wd:nn { \l__box_right_new_dim }

```

```

28194         {
28195             \tex_kern:D \l__box_right_new_dim
28196             \box_use_drop:N \l__box_internal_box
28197             \tex_hss:D
28198         }
28199     }
28200     {
28201         \box_set_wd:Nn \l__box_internal_box { \l__box_right_new_dim }
28202         \hbox:n
28203         {
28204             \tex_kern:D \c_zero_dim
28205             \box_use_drop:N \l__box_internal_box
28206             \tex_hss:D
28207         }
28208     }
28209 }

```

(End definition for `__box_resize_common:N`.)

```
28210 \</package>
```

43 l3coffins Implementation

```

28211 \*package>
28212 \@@=coffin>

```

43.1 Coffins: data structures and general variables

`\l__coffin_internal_box` Scratch variables.

```

\l__coffin_internal_dim 28213 \box_new:N \l__coffin_internal_box
\l__coffin_internal_tl 28214 \dim_new:N \l__coffin_internal_dim
28215 \tl_new:N \l__coffin_internal_tl

```

(End definition for `\l__coffin_internal_box`, `\l__coffin_internal_dim`, and `\l__coffin_internal_tl`.)

`\c__coffin_corners_prop` The “corners”; of a coffin define the real content, as opposed to the $\mathrm{T}_{\mathrm{E}}\mathrm{X}$ bounding box. They all start off in the same place, of course.

```

28216 \prop_const_from_keyval:Nn \c__coffin_corners_prop
28217 {
28218     tl = { Opt } { Opt } ,
28219     tr = { Opt } { Opt } ,
28220     bl = { Opt } { Opt } ,
28221     br = { Opt } { Opt } ,
28222 }

```

(End definition for `\c__coffin_corners_prop`.)

`\c__coffin_poles_prop` Pole positions are given for horizontal, vertical and reference-point based values.

```

28223 \prop_const_from_keyval:Nn \c__coffin_poles_prop
28224 {
28225     l  = { Opt } { Opt } { Opt } { 1000pt } ,
28226     hc = { Opt } { Opt } { Opt } { 1000pt } ,
28227     r  = { Opt } { Opt } { Opt } { 1000pt } ,

```

```

28228      b = { 0pt } { 0pt } { 1000pt } { 0pt } ,
28229      vc = { 0pt } { 0pt } { 1000pt } { 0pt } ,
28230      t  = { 0pt } { 0pt } { 1000pt } { 0pt } ,
28231      B  = { 0pt } { 0pt } { 1000pt } { 0pt } ,
28232      H  = { 0pt } { 0pt } { 1000pt } { 0pt } ,
28233      T  = { 0pt } { 0pt } { 1000pt } { 0pt } ,
28234      }

```

(End definition for \c__coffin_poles_prop.)

\l__coffin_slope_A_fp Used for calculations of intersections.

```

\l__coffin_slope_B_fp
28235 \fp_new:N \l__coffin_slope_A_fp
28236 \fp_new:N \l__coffin_slope_B_fp

```

(End definition for \l__coffin_slope_A_fp and \l__coffin_slope_B_fp.)

\l__coffin_error_bool For propagating errors so that parts of the code can work around them.

```

28237 \bool_new:N \l__coffin_error_bool

```

(End definition for \l__coffin_error_bool.)

\l__coffin_offset_x_dim The offset between two sets of coffin handles when typesetting. These values are corrected from those requested in an alignment for the positions of the handles.

```

28238 \dim_new:N \l__coffin_offset_x_dim
28239 \dim_new:N \l__coffin_offset_y_dim

```

(End definition for \l__coffin_offset_x_dim and \l__coffin_offset_y_dim.)

\l__coffin_pole_a_tl Needed for finding the intersection of two poles.

```

\l__coffin_pole_b_tl
28240 \tl_new:N \l__coffin_pole_a_tl
28241 \tl_new:N \l__coffin_pole_b_tl

```

(End definition for \l__coffin_pole_a_tl and \l__coffin_pole_b_tl.)

\l__coffin_x_dim For calculating intersections and so forth.

```

\l__coffin_y_dim
28242 \dim_new:N \l__coffin_x_dim
28243 \dim_new:N \l__coffin_y_dim
\l__coffin_x_prime_dim
28244 \dim_new:N \l__coffin_x_prime_dim
\l__coffin_y_prime_dim
28245 \dim_new:N \l__coffin_y_prime_dim

```

(End definition for \l__coffin_x_dim and others.)

43.2 Basic coffin functions

There are a number of basic functions needed for creating coffins and placing material in them. This all relies on the following data structures.

__coffin_to_value:N Coffins are a two-part structure and we rely on the internal nature of box allocation to make everything work. As such, we need an interface to turn coffin identifiers into numbers. For the purposes here, the signature allowed is N despite the nature of the underlying primitive.

```

28246 \cs_new_eq:NN \__coffin_to_value:N \tex_number:D

```

(End definition for __coffin_to_value:N.)

`\coffin_if_exist_p:N` Several of the higher-level coffin functions would give multiple errors if the coffin does not exist. A cleaner way to handle this is provided here: both the box and the coffin structure are checked.

```

\coffin_if_exist_p:c
\coffin_if_exist:NTF
\coffin_if_exist:cTF
28247 \prg_new_conditional:Npnn \coffin_if_exist:N #1 { p , T , F , TF }
28248 {
28249   \cs_if_exist:NTF #1
28250   {
28251     \cs_if_exist:cTF { coffin ~ \__coffin_to_value:N #1 ~ poles }
28252     { \prg_return_true: }
28253     { \prg_return_false: }
28254   }
28255   { \prg_return_false: }
28256 }
28257 \prg_generate_conditional_variant:Nnn \coffin_if_exist:N
28258 { c } { p , T , F , TF }

```

(End definition for `\coffin_if_exist:NTF`. This function is documented on page 253.)

`__coffin_if_exist:NT` Several of the higher-level coffin functions would give multiple errors if the coffin does not exist. So a wrapper is provided to deal with this correctly, issuing an error on erroneous use.

```

28259 \cs_new_protected:Npn \__coffin_if_exist:NT #1#2
28260 {
28261   \coffin_if_exist:NTF #1
28262   { #2 }
28263   {
28264     \_kernel_msg_error:nxx { kernel } { unknown-coffin }
28265     { \token_to_str:N #1 }
28266   }
28267 }

```

(End definition for `__coffin_if_exist:NT`.)

`\coffin_clear:N` Clearing coffins means emptying the box and resetting all of the structures.

```

\coffin_clear:c
\coffin_gclear:N
\coffin_gclear:c
28268 \cs_new_protected:Npn \coffin_clear:N #1
28269 {
28270   \__coffin_if_exist:NT #1
28271   {
28272     \box_clear:N #1
28273     \__coffin_reset_structure:N #1
28274   }
28275 }
28276 \cs_generate_variant:Nn \coffin_clear:N { c }
28277 \cs_new_protected:Npn \coffin_gclear:N #1
28278 {
28279   \__coffin_if_exist:NT #1
28280   {
28281     \box_gclear:N #1
28282     \__coffin_greset_structure:N #1
28283   }
28284 }
28285 \cs_generate_variant:Nn \coffin_gclear:N { c }

```

(End definition for `\coffin_clear:N` and `\coffin_gclear:N`. These functions are documented on page 253.)

\coffin_new:N Creating a new coffin means making the underlying box and adding the data structures. The **\debug_suspend:** and **\debug_resume:** functions prevent **\prop_gclear_new:c** from writing useless information to the log file.

```

28286 \cs_new_protected:Npn \coffin_new:N #1
28287 {
28288   \box_new:N #1
28289   \debug_suspend:
28290   \prop_gclear_new:c { coffin ~ \__coffin_to_value:N #1 ~ corners }
28291   \prop_gclear_new:c { coffin ~ \__coffin_to_value:N #1 ~ poles }
28292   \prop_gset_eq:cN { coffin ~ \__coffin_to_value:N #1 ~ corners }
28293     \c__coffin_corners_prop
28294   \prop_gset_eq:cN { coffin ~ \__coffin_to_value:N #1 ~ poles }
28295     \c__coffin_poles_prop
28296   \debug_resume:
28297 }
28298 \cs_generate_variant:Nn \coffin_new:N { c }

```

(End definition for **\coffin_new:N**. This function is documented on page 253.)

\hcoffin_set:Nn Horizontal coffins are relatively easy: set the appropriate box, reset the structures then
\hcoffin_set:cn update the handle positions.

```

28299 \cs_new_protected:Npn \hcoffin_set:Nn #1#2
28300 {
28301   \__coffin_if_exist:NT #1
28302   {
28303     \hbox_set:Nn #1
28304     {
28305       \color_ensure_current:
28306       #2
28307     }
28308     \__coffin_update:N #1
28309   }
28310 }
28311 \cs_generate_variant:Nn \hcoffin_set:Nn { c }
28312 \cs_new_protected:Npn \hcoffin_gset:Nn #1#2
28313 {
28314   \__coffin_if_exist:NT #1
28315   {
28316     \hbox_gset:Nn #1
28317     {
28318       \color_ensure_current:
28319       #2
28320     }
28321     \__coffin_gupdate:N #1
28322   }
28323 }
28324 \cs_generate_variant:Nn \hcoffin_gset:Nn { c }

```

(End definition for **\hcoffin_set:Nn** and **\hcoffin_gset:Nn**. These functions are documented on page 253.)

\vcoffin_set:Nnn Setting vertical coffins is more complex. First, the material is typeset with a given width.
\vcoffin_set:cnn The default handles and poles are set as for a horizontal coffin, before finding the top
\vcoffin_gset:Nnn baseline using a temporary box. No **\color_ensure_current:** here as that would add a
\vcoffin_gset:cnn

```

\__coffin_set_vertical:NnnNN
\__coffin_set_vertical_aux:

```

whatsit to the start of the vertical box and mess up the location of the T pole (see *TeX by Topic* for discussion of the `\vtop` primitive, used to do the measuring).

```

28325 \cs_new_protected:Npn \vcoffin_set:Nnn #1#2#3
28326 {
28327   \__coffin_set_vertical:NnnNN #1 {#2} {#3}
28328   \vbox_set:Nn \__coffin_update:N
28329 }
28330 \cs_generate_variant:Nn \vcoffin_set:Nnn { c }
28331 \cs_new_protected:Npn \vcoffin_gset:Nnn #1#2#3
28332 {
28333   \__coffin_set_vertical:NnnNN #1 {#2} {#3}
28334   \vbox_gset:Nn \__coffin_gupdate:N
28335 }
28336 \cs_generate_variant:Nn \vcoffin_gset:Nnn { c }
28337 \cs_new_protected:Npn \__coffin_set_vertical:NnnNN #1#2#3#4#5
28338 {
28339   \__coffin_if_exist:NT #1
28340   {
28341     #4 #1
28342     {
28343       \dim_set:Nn \tex_hsize:D {#2}
28344       \__coffin_set_vertical_aux:
28345       #3
28346     }
28347     #5 #1
28348     \vbox_set_top:Nn \l__coffin_internal_box { \vbox_unpack:N #1 }
28349     \__coffin_set_pole:Nnx #1 { T }
28350     {
28351       { Opt }
28352       {
28353         \dim_eval:n
28354         { \box_ht:N #1 - \box_ht:N \l__coffin_internal_box }
28355       }
28356       { 1000pt }
28357       { Opt }
28358     }
28359     \box_clear:N \l__coffin_internal_box
28360   }
28361 }
28362 \cs_new_protected:Npx \__coffin_set_vertical_aux:
28363 {
28364   \bool_lazy_and:nnT
28365   { \cs_if_exist_p:N \fmtname }
28366   { \str_if_eq_p:Vn \fmtname { LaTeX2e } }
28367   {
28368     \dim_set_eq:NN \exp_not:N \linewidth \tex_hsize:D
28369     \dim_set_eq:NN \exp_not:N \columnwidth \tex_hsize:D
28370   }
28371 }

```

(End definition for `\vcoffin_set:Nnn` and others. These functions are documented on page 254.)

`\hcoffin_set:Nw` These are the “begin”/“end” versions of the above: watch the grouping!

`\hcoffin_set:cw` 28372 `\cs_new_protected:Npn \hcoffin_set:Nw #1`

`\hcoffin_gset:Nw`

`\hcoffin_gset:cw`

`\hcoffin_set_end:`

`\hcoffin_gset_end:`

```

28373 {
28374   \__coffin_if_exist:NT #1
28375   {
28376     \hbox_set:Nw #1 \color_ensure_current:
28377     \cs_set_protected:Npn \hcoffin_set_end:
28378     {
28379       \hbox_set_end:
28380       \__coffin_update:N #1
28381     }
28382   }
28383 }
28384 \cs_generate_variant:Nn \hcoffin_set:Nw { c }
28385 \cs_new_protected:Npn \hcoffin_gset:Nw #1
28386 {
28387   \__coffin_if_exist:NT #1
28388   {
28389     \hbox_gset:Nw #1 \color_ensure_current:
28390     \cs_set_protected:Npn \hcoffin_gset_end:
28391     {
28392       \hbox_gset_end:
28393       \__coffin_gupdate:N #1
28394     }
28395   }
28396 }
28397 \cs_generate_variant:Nn \hcoffin_gset:Nw { c }
28398 \cs_new_protected:Npn \hcoffin_set_end: { }
28399 \cs_new_protected:Npn \hcoffin_gset_end: { }

```

(End definition for \hcoffin_set:Nw and others. These functions are documented on page 254.)

```

\vcoffin_set:Nnw The same for vertical coffins.
\vcoffin_set:cnw
\vcoffin_gset:Nnw
\vcoffin_gset:cnw
\__coffin_set_vertical:NnNNNw
\vcoffin_set_end:
\vcoffin_gset_end:
28400 \cs_new_protected:Npn \vcoffin_set:Nnw #1#2
28401 {
28402   \__coffin_set_vertical:NnNNNw #1 {#2} \vbox_set:Nw
28403   \vcoffin_set_end:
28404   \vbox_set_end: \__coffin_update:N
28405 }
28406 \cs_generate_variant:Nn \vcoffin_set:Nnw { c }
28407 \cs_new_protected:Npn \vcoffin_gset:Nnw #1#2
28408 {
28409   \__coffin_set_vertical:NnNNNw #1 {#2} \vbox_gset:Nw
28410   \vcoffin_gset_end:
28411   \vbox_gset_end: \__coffin_gupdate:N
28412 }
28413 \cs_generate_variant:Nn \vcoffin_gset:Nnw { c }
28414 \cs_new_protected:Npn \__coffin_set_vertical:NnNNNw #1#2#3#4#5#6
28415 {
28416   \__coffin_if_exist:NT #1
28417   {
28418     #3 #1
28419     \dim_set:Nn \tex_hsize:D {#2}
28420     \__coffin_set_vertical_aux:
28421     \cs_set_protected:Npn #4
28422     {

```



```

28423         #5
28424         #6 #1
28425         \vbox_set_top:Nn \l__coffin_internal_box { \vbox_unpack:N #1 }
28426         \__coffin_set_pole:Nnx #1 { T }
28427         {
28428             { Opt }
28429             {
28430                 \dim_eval:n
28431                 { \box_ht:N #1 - \box_ht:N \l__coffin_internal_box }
28432             }
28433             { 1000pt }
28434             { Opt }
28435         }
28436         \box_clear:N \l__coffin_internal_box
28437     }
28438 }
28439 }
28440 \cs_new_protected:Npn \vcoffin_set_end: { }
28441 \cs_new_protected:Npn \vcoffin_gset_end: { }

```

(End definition for \vcoffin_set:Nnw and others. These functions are documented on page 254.)

```

\coffin_set_eq:NN Setting two coffins equal is just a wrapper around other functions.
\coffin_set_eq:Nc 28442 \cs_new_protected:Npn \coffin_set_eq:NN #1#2
\coffin_set_eq:cN 28443 {
\coffin_set_eq:cc 28444     \__coffin_if_exist:NT #1
\coffin_gset_eq:NN 28445     {
\coffin_gset_eq:Nc 28446         \box_set_eq:NN #1 #2
\coffin_gset_eq:cN 28447         \prop_set_eq:cc { coffin ~ \__coffin_to_value:N #1 ~ corners }
\coffin_gset_eq:cc 28448         { coffin ~ \__coffin_to_value:N #2 ~ corners }
28449         \prop_set_eq:cc { coffin ~ \__coffin_to_value:N #1 ~ poles }
28450         { coffin ~ \__coffin_to_value:N #2 ~ poles }
28451     }
28452 }
28453 \cs_generate_variant:Nn \coffin_set_eq:NN { c , Nc , cc }
28454 \cs_new_protected:Npn \coffin_gset_eq:NN #1#2
28455 {
28456     \__coffin_if_exist:NT #1
28457     {
28458         \box_gset_eq:NN #1 #2
28459         \prop_gset_eq:cc { coffin ~ \__coffin_to_value:N #1 ~ corners }
28460         { coffin ~ \__coffin_to_value:N #2 ~ corners }
28461         \prop_gset_eq:cc { coffin ~ \__coffin_to_value:N #1 ~ poles }
28462         { coffin ~ \__coffin_to_value:N #2 ~ poles }
28463     }
28464 }
28465 \cs_generate_variant:Nn \coffin_gset_eq:NN { c , Nc , cc }

```

(End definition for \coffin_set_eq:NN and \coffin_gset_eq:NN. These functions are documented on page 253.)

\c_empty_coffin Special coffins: these cannot be set up earlier as they need \coffin_new:N. The empty coffin is set as a box as the full coffin-setting system needs some material which is not yet available. The empty coffin is created entirely by hand: not everything is in place yet.

```

\l__coffin_aligned_coffin
\l__coffin_aligned_internal_coffin

```

```

28466 \coffin_new:N \c_empty_coffin
28467 \coffin_new:N \l__coffin_aligned_coffin
28468 \coffin_new:N \l__coffin_aligned_internal_coffin

```

(End definition for `\c_empty_coffin`, `\l__coffin_aligned_coffin`, and `\l__coffin_aligned_internal_coffin`. This variable is documented on page 257.)

```

\l_tmpa_coffin The usual scratch space.
\l_tmpb_coffin
\g_tmpa_coffin
\g_tmpb_coffin

```

(End definition for `\l_tmpa_coffin` and others. These variables are documented on page 257.)

43.3 Measuring coffins

```

\coffin_dp:N Coffins are just boxes when it comes to measurement. However, semantically a separate
\coffin_dp:c set of functions are required.
\coffin_ht:N
\coffin_ht:c
\coffin_wd:N
\coffin_wd:c

```

(End definition for `\coffin_dp:N`, `\coffin_ht:N`, and `\coffin_wd:N`. These functions are documented on page 256.)

43.4 Coffins: handle and pole management

```

\__coffin_get_pole:NnN A simple wrapper around the recovery of a coffin pole, with some error checking and
recovery built-in.

```

```

28479 \cs_new_protected:Npn \__coffin_get_pole:NnN #1#2#3
28480 {
28481   \prop_get:cnNF
28482   { coffin ~ \__coffin_to_value:N #1 ~ poles } {#2} #3
28483   {
28484     \__kernel_msg_error:nxxx { kernel } { unknown-coffin-pole }
28485     { \exp_not:n {#2} } { \token_to_str:N #1 }
28486     \tl_set:Nn #3 { { 0pt } { 0pt } { 0pt } { 0pt } }
28487   }
28488 }

```

(End definition for `__coffin_get_pole:NnN`.)

```

\__coffin_reset_structure:N Resetting the structure is a simple copy job.

```

```

\__coffin_greset_structure:N
28489 \cs_new_protected:Npn \__coffin_reset_structure:N #1
28490 {
28491   \prop_set_eq:cN { coffin ~ \__coffin_to_value:N #1 ~ corners }
28492   \c__coffin_corners_prop
28493   \prop_set_eq:cN { coffin ~ \__coffin_to_value:N #1 ~ poles }
28494   \c__coffin_poles_prop
28495 }

```

```

28496 \cs_new_protected:Npn \__coffin_greset_structure:N #1
28497 {
28498   \prop_gset_eq:cN { coffin ~ \__coffin_to_value:N #1 ~ corners }
28499   \c__coffin_corners_prop
28500   \prop_gset_eq:cN { coffin ~ \__coffin_to_value:N #1 ~ poles }
28501   \c__coffin_poles_prop
28502 }

```

(End definition for __coffin_reset_structure:N and __coffin_greset_structure:N.)

\coffin_set_horizontal_pole:Nnn
\coffin_set_horizontal_pole:cnm
\coffin_gset_horizontal_pole:Nnn
\coffin_gset_horizontal_pole:cnm
__coffin_set_horizontal_pole:NnnN
\coffin_set_vertical_pole:Nnn
\coffin_set_vertical_pole:cnm
\coffin_gset_vertical_pole:Nnn
\coffin_gset_vertical_pole:cnm
__coffin_set_vertical_pole:NnnN
__coffin_set_pole:Nnn
__coffin_set_pole:Nnx

Setting the pole of a coffin at the user/designer level requires a bit more care. The idea here is to provide a reasonable interface to the system, then to do the setting with full expansion. The three-argument version is used internally to do a direct setting.

```

28503 \cs_new_protected:Npn \coffin_set_horizontal_pole:Nnn #1#2#3
28504 { \__coffin_set_horizontal_pole:NnnN #1 {#2} {#3} \prop_put:cnx }
28505 \cs_generate_variant:Nn \coffin_set_horizontal_pole:Nnn { c }
28506 \cs_new_protected:Npn \coffin_gset_horizontal_pole:Nnn #1#2#3
28507 { \__coffin_set_horizontal_pole:NnnN #1 {#2} {#3} \prop_gput:cnx }
28508 \cs_generate_variant:Nn \coffin_gset_horizontal_pole:Nnn { c }
28509 \cs_new_protected:Npn \__coffin_set_horizontal_pole:NnnN #1#2#3#4
28510 {
28511   \__coffin_if_exist:NT #1
28512   {
28513     #4 { coffin ~ \__coffin_to_value:N #1 ~ poles }
28514     {#2}
28515     {
28516       { Opt } { \dim_eval:n {#3} }
28517       { 1000pt } { Opt }
28518     }
28519   }
28520 }
28521 \cs_new_protected:Npn \coffin_set_vertical_pole:Nnn #1#2#3
28522 { \__coffin_set_vertical_pole:NnnN #1 {#2} {#3} \prop_put:cnx }
28523 \cs_generate_variant:Nn \coffin_set_vertical_pole:Nnn { c }
28524 \cs_new_protected:Npn \coffin_gset_vertical_pole:Nnn #1#2#3
28525 { \__coffin_set_vertical_pole:NnnN #1 {#2} {#3} \prop_gput:cnx }
28526 \cs_generate_variant:Nn \coffin_gset_vertical_pole:Nnn { c }
28527 \cs_new_protected:Npn \__coffin_set_vertical_pole:NnnN #1#2#3#4
28528 {
28529   \__coffin_if_exist:NT #1
28530   {
28531     #4 { coffin ~ \__coffin_to_value:N #1 ~ poles }
28532     {#2}
28533     {
28534       { \dim_eval:n {#3} } { Opt }
28535       { Opt } { 1000pt }
28536     }
28537   }
28538 }
28539 \cs_new_protected:Npn \__coffin_set_pole:Nnn #1#2#3
28540 {
28541   \prop_put:cnm { coffin ~ \__coffin_to_value:N #1 ~ poles }
28542   {#2} {#3}
28543 }

```

```
28544 \cs_generate_variant:Nn \__coffin_set_pole:Nnn { Nnx }
```

(End definition for `\coffin_set_horizontal_pole:Nnn` and others. These functions are documented on page 254.)

```
\__coffin_update:N
```

Simple shortcuts.

```
\__coffin_gupdate:N
```

```
28545 \cs_new_protected:Npn \__coffin_update:N #1
28546 {
28547   \__coffin_reset_structure:N #1
28548   \__coffin_update_corners:N #1
28549   \__coffin_update_poles:N #1
28550 }
28551 \cs_new_protected:Npn \__coffin_gupdate:N #1
28552 {
28553   \__coffin_greset_structure:N #1
28554   \__coffin_gupdate_corners:N #1
28555   \__coffin_gupdate_poles:N #1
28556 }
```

(End definition for `__coffin_update:N` and `__coffin_gupdate:N`.)

```
\__coffin_update_corners:N
```

Updating the corners of a coffin is straight-forward as at this stage there can be no rotation. So the corners of the content are just those of the underlying TeX box.

```
\__coffin_gupdate_corners:N
```

```
\__coffin_update_corners:NN
```

```
\__coffin_update_corners:NNN
```

```
28557 \cs_new_protected:Npn \__coffin_update_corners:N #1
28558 { \__coffin_update_corners:NN #1 \prop_put:Nnx }
28559 \cs_new_protected:Npn \__coffin_gupdate_corners:N #1
28560 { \__coffin_update_corners:NN #1 \prop_gput:Nnx }
28561 \cs_new_protected:Npn \__coffin_update_corners:NN #1#2
28562 {
28563   \exp_args:Nc \__coffin_update_corners:NNN
28564   { coffin ~ \__coffin_to_value:N #1 ~ corners }
28565   #1 #2
28566 }
28567 \cs_new_protected:Npn \__coffin_update_corners:NNN #1#2#3
28568 {
28569   #3 #1
28570   { tl }
28571   { { Opt } { \dim_eval:n { \box_ht:N #2 } } }
28572   #3 #1
28573   { tr }
28574   {
28575     { \dim_eval:n { \box_wd:N #2 } }
28576     { \dim_eval:n { \box_ht:N #2 } }
28577   }
28578   #3 #1
28579   { bl }
28580   { { Opt } { \dim_eval:n { -\box_dp:N #2 } } }
28581   #3 #1
28582   { br }
28583   {
28584     { \dim_eval:n { \box_wd:N #2 } }
28585     { \dim_eval:n { -\box_dp:N #2 } }
28586   }
28587 }
```

(End definition for `_coffin_update_corners:N` and others.)

`_coffin_update_poles:N`
`_coffin_gupdate_poles:N`
`_coffin_update_poles:NN`
`_coffin_update_poles:NNN`

This function is called when a coffin is set, and updates the poles to reflect the nature of size of the box. Thus this function only alters poles where the default position is dependent on the size of the box. It also does not set poles which are relevant only to vertical coffins.

```

28588 \cs_new_protected:Npn \_coffin_update_poles:N #1
28589   { \_coffin_update_poles:NN #1 \prop_put:Nnx }
28590 \cs_new_protected:Npn \_coffin_gupdate_poles:N #1
28591   { \_coffin_update_poles:NN #1 \prop_gput:Nnx }
28592 \cs_new_protected:Npn \_coffin_update_poles:NN #1#2
28593   {
28594     \exp_args:Nc \_coffin_update_poles:NNN
28595       { coffin ~ \_coffin_to_value:N #1 ~ poles }
28596     #1 #2
28597   }
28598 \cs_new_protected:Npn \_coffin_update_poles:NNN #1#2#3
28599   {
28600     #3 #1 { hc }
28601     {
28602       { \dim_eval:n { 0.5 \box_wd:N #2 } }
28603       { 0pt } { 0pt } { 1000pt }
28604     }
28605     #3 #1 { r }
28606     {
28607       { \dim_eval:n { \box_wd:N #2 } }
28608       { 0pt } { 0pt } { 1000pt }
28609     }
28610     #3 #1 { vc }
28611     {
28612       { 0pt }
28613       { \dim_eval:n { ( \box_ht:N #2 - \box_dp:N #2 ) / 2 } }
28614       { 1000pt }
28615       { 0pt }
28616     }
28617     #3 #1 { t }
28618     {
28619       { 0pt }
28620       { \dim_eval:n { \box_ht:N #2 } }
28621       { 1000pt }
28622       { 0pt }
28623     }
28624     #3 #1 { b }
28625     {
28626       { 0pt }
28627       { \dim_eval:n { -\box_dp:N #2 } }
28628       { 1000pt }
28629       { 0pt }
28630     }
28631   }

```

(End definition for `_coffin_update_poles:N` and others.)

43.5 Coffins: calculation of pole intersections

The lead off in finding intersections is to recover the two poles and then hand off to the auxiliary for the actual calculation. There may of course not be an intersection, for which an error trap is needed.

```

28632 \cs_new_protected:Npn \__coffin_calculate_intersection:Nnn #1#2#3
28633 {
28634   \__coffin_get_pole:NnN #1 {#2} \l__coffin_pole_a_tl
28635   \__coffin_get_pole:NnN #1 {#3} \l__coffin_pole_b_tl
28636   \bool_set_false:N \l__coffin_error_bool
28637   \exp_last_two_unbraced:Noo
28638   \__coffin_calculate_intersection:nnnnnnnn
28639   \l__coffin_pole_a_tl \l__coffin_pole_b_tl
28640   \bool_if:NT \l__coffin_error_bool
28641   {
28642     \kernel_msg_error:nn { kernel } { no-pole-intersection }
28643     \dim_zero:N \l__coffin_x_dim
28644     \dim_zero:N \l__coffin_y_dim
28645   }
28646 }
```

The two poles passed here each have four values (as dimensions), (a, b, c, d) and (a', b', c', d') . These are arguments 1–4 and 5–8, respectively. In both cases a and b are the co-ordinates of a point on the pole and c and d define the direction of the pole. Finding the intersection depends on the directions of the poles, which are given by d/c and d'/c' . However, if one of the poles is either horizontal or vertical then one or more of c, d, c' and d' are zero and a special case is needed.

```

28647 \cs_new_protected:Npn \__coffin_calculate_intersection:nnnnnnnn
28648   #1#2#3#4#5#6#7#8
28649 {
28650   \dim_compare:nNnTF {#3} = \c_zero_dim
```

The case where the first pole is vertical. So the x -component of the interaction is at a . There is then a test on the second pole: if it is also vertical then there is an error.

```

28651   {
28652     \dim_set:Nn \l__coffin_x_dim {#1}
28653     \dim_compare:nNnTF {#7} = \c_zero_dim
28654     { \bool_set_true:N \l__coffin_error_bool }
```

The second pole may still be horizontal, in which case the y -component of the intersection is b' . If not,

$$y = \frac{d'}{c'} (a - a') + b'$$

with the x -component already known to be $\#1$.

```

28655   {
28656     \dim_set:Nn \l__coffin_y_dim
28657     {
28658       \dim_compare:nNnTF {#8} = \c_zero_dim
28659       {#6}
28660       {
28661         \fp_to_dim:n
28662         {
28663           ( \dim_to_fp:n {#8} / \dim_to_fp:n {#7} )
28664           * ( \dim_to_fp:n {#1} - \dim_to_fp:n {#5} )
```

```

28665         + \dim_to_fp:n {#6}
28666     }
28667 }
28668 }
28669 }
28670 }

```

If the first pole is not vertical then it may be horizontal. If so, then the procedure is essentially the same as that already done but with the x - and y -components interchanged.

```

28671 {
28672     \dim_compare:nNnTF {#4} = \c_zero_dim
28673     {
28674         \dim_set:Nn \l__coffin_y_dim {#2}
28675         \dim_compare:nNnTF {#8} = { \c_zero_dim }
28676         { \bool_set_true:N \l__coffin_error_bool }
28677     }

```

Now we deal with the case where the second pole may be vertical, or if not we have

$$x = \frac{c'}{d'} (b - b') + a'$$

which is again handled by the same auxiliary.

```

28678     \dim_set:Nn \l__coffin_x_dim
28679     {
28680         \dim_compare:nNnTF {#7} = \c_zero_dim
28681         {#5}
28682         {
28683             \fp_to_dim:n
28684             {
28685                 ( \dim_to_fp:n {#7} / \dim_to_fp:n {#8} )
28686                 * ( \dim_to_fp:n {#4} - \dim_to_fp:n {#6} )
28687                 + \dim_to_fp:n {#5}
28688             }
28689         }
28690     }
28691 }
28692 }

```

The first pole is neither horizontal nor vertical. To avoid even more complexity, we now work out both slopes and pass to an auxiliary.

```

28693 {
28694     \use:x
28695     {
28696         \__coffin_calculate_intersection:nnnnnn
28697         { \dim_to_fp:n {#4} / \dim_to_fp:n {#3} }
28698         { \dim_to_fp:n {#8} / \dim_to_fp:n {#7} }
28699     }
28700     {#1} {#2} {#5} {#6}
28701 }
28702 }
28703 }

```

Assuming the two poles are not parallel, then the intersection point is found in two steps. First we find the x -value with

$$x = \frac{sa - s'a' - b + b'}{s - s'}$$

and then finding the y -value with

$$y = s(x - a) + b$$

```

28704 \cs_set_protected:Npn \__coffin_calculate_intersection:nnnnnn #1#2#3#4#5#6
28705 {
28706   \fp_compare:nNnTF {#1} = {#2}
28707   { \bool_set_true:N \l__coffin_error_bool }
28708   {
28709     \dim_set:Nn \l__coffin_x_dim
28710     {
28711       \fp_to_dim:n
28712       {
28713         (
28714           #1 * \dim_to_fp:n {#3}
28715           - #2 * \dim_to_fp:n {#5}
28716           - \dim_to_fp:n {#4}
28717           + \dim_to_fp:n {#6}
28718         )
28719         /
28720         ( #1 - #2 )
28721       }
28722     }
28723     \dim_set:Nn \l__coffin_y_dim
28724     {
28725       \fp_to_dim:n
28726       {
28727         #1 * ( \l__coffin_x_dim - \dim_to_fp:n {#3} )
28728         + \dim_to_fp:n {#4}
28729       }
28730     }
28731   }
28732 }

```

(End definition for `__coffin_calculate_intersection:Nnn`, `__coffin_calculate_intersection:nnnnnnnn`, and `__coffin_calculate_intersection:nnnnnn`.)

43.6 Affine transformations

`\l__coffin_sin_fp` Used for rotations to get the sine and cosine values.

```

\l__coffin_cos_fp 28733 \fp_new:N \l__coffin_sin_fp
28734 \fp_new:N \l__coffin_cos_fp

```

(End definition for `\l__coffin_sin_fp` and `\l__coffin_cos_fp`.)

`\l__coffin_bounding_prop` A property list for the bounding box of a coffin. This is only needed during the rotation, so there is just the one.

```
28735 \prop_new:N \l__coffin_bounding_prop
```

(End definition for `\l__coffin_bounding_prop`.)

`\l__coffin_corners_prop` Used to avoid needing to track scope for intermediate steps.

```

\l__coffin_poles_prop 28736 \prop_new:N \l__coffin_corners_prop
28737 \prop_new:N \l__coffin_poles_prop

```


(End definition for \l__coffin_corners_prop and \l__coffin_poles_prop.)

\l__coffin_bounding_shift_dim The shift of the bounding box of a coffin from the real content.

28738 \dim_new:N \l__coffin_bounding_shift_dim

(End definition for \l__coffin_bounding_shift_dim.)

\l__coffin_left_corner_dim These are used to hold maxima for the various corner values: these thus define the minimum size of the bounding box after rotation.

\l__coffin_right_corner_dim

\l__coffin_bottom_corner_dim

\l__coffin_top_corner_dim

28739 \dim_new:N \l__coffin_left_corner_dim

28740 \dim_new:N \l__coffin_right_corner_dim

28741 \dim_new:N \l__coffin_bottom_corner_dim

28742 \dim_new:N \l__coffin_top_corner_dim

(End definition for \l__coffin_left_corner_dim and others.)

\coffin_rotate:Nn Rotating a coffin requires several steps which can be conveniently run together. The sine and cosine of the angle in degrees are computed. This is then used to set \l__coffin_sin_fp and \l__coffin_cos_fp, which are carried through unchanged for the rest of the procedure.

\coffin_rotate:cn

\coffin_grotate:Nn

\coffin_grotate:cn

__coffin_rotate:NnNNN

28743 \cs_new_protected:Npn \coffin_rotate:Nn #1#2

28744 { __coffin_rotate:NnNNN #1 {#2} \box_rotate:Nn \prop_set_eq:cn \hbox_set:Nn }

28745 \cs_generate_variant:Nn \coffin_rotate:Nn { c }

28746 \cs_new_protected:Npn \coffin_grotate:Nn #1#2

28747 { __coffin_rotate:NnNNN #1 {#2} \box_grotate:Nn \prop_gset_eq:cn \hbox_gset:Nn }

28748 \cs_generate_variant:Nn \coffin_grotate:Nn { c }

28749 \cs_new_protected:Npn __coffin_rotate:NnNNN #1#2#3#4#5

28750 {

28751 \fp_set:Nn \l__coffin_sin_fp { sind (#2) }

28752 \fp_set:Nn \l__coffin_cos_fp { cosd (#2) }

Use a local copy of the property lists to avoid needing to pass the name and scope around.

28753 \prop_set_eq:Nc \l__coffin_corners_prop

28754 { coffin ~ __coffin_to_value:N #1 ~ corners }

28755 \prop_set_eq:Nc \l__coffin_poles_prop

28756 { coffin ~ __coffin_to_value:N #1 ~ poles }

The corners and poles of the coffin can now be rotated around the origin. This is best achieved using mapping functions.

28757 \prop_map_inline:Nn \l__coffin_corners_prop

28758 { __coffin_rotate_corner:Nnnn #1 {##1} ##2 }

28759 \prop_map_inline:Nn \l__coffin_poles_prop

28760 { __coffin_rotate_pole:Nnnnnn #1 {##1} ##2 }

The bounding box of the coffin needs to be rotated, and to do this the corners have to be found first. They are then rotated in the same way as the corners of the coffin material itself.

28761 __coffin_set_bounding:N #1

28762 \prop_map_inline:Nn \l__coffin_bounding_prop

28763 { __coffin_rotate_bounding:nnn {##1} ##2 }

At this stage, there needs to be a calculation to find where the corners of the content and the box itself will end up.

28764 __coffin_find_corner_maxima:N #1

28765 __coffin_find_bounding_shift:

28766 #3 #1 {#2}

The correction of the box position itself takes place here. The idea is that the bounding box for a coffin is tight up to the content, and has the reference point at the bottom-left. The x -direction is handled by moving the content by the difference in the positions of the bounding box and the content left edge. The y -direction is dealt with by moving the box down by any depth it has acquired. The internal box is used here to allow for the next step.

```

28767 \hbox_set:Nn \l__coffin_internal_box
28768 {
28769   \tex_kern:D
28770   \dim_eval:n
28771     { \l__coffin_bounding_shift_dim - \l__coffin_left_corner_dim }
28772   \exp_stop_f:
28773   \box_move_down:nn { \l__coffin_bottom_corner_dim }
28774   { \box_use:N #1 }
28775 }

```

If there have been any previous rotations then the size of the bounding box will be bigger than the contents. This can be corrected easily by setting the size of the box to the height and width of the content. As this operation requires setting box dimensions and these transcend grouping, the safe way to do this is to use the internal box and to reset the result into the target box.

```

28776 \box_set_ht:Nn \l__coffin_internal_box
28777   { \l__coffin_top_corner_dim - \l__coffin_bottom_corner_dim }
28778 \box_set_dp:Nn \l__coffin_internal_box { 0pt }
28779 \box_set_wd:Nn \l__coffin_internal_box
28780   { \l__coffin_right_corner_dim - \l__coffin_left_corner_dim }
28781 #5 #1 { \box_use_drop:N \l__coffin_internal_box }

```

The final task is to move the poles and corners such that they are back in alignment with the box reference point.

```

28782 \prop_map_inline:Nn \l__coffin_corners_prop
28783   { \__coffin_shift_corner:Nnnn #1 {##1} ##2 }
28784 \prop_map_inline:Nn \l__coffin_poles_prop
28785   { \__coffin_shift_pole:Nnnnnn #1 {##1} ##2 }

```

Update the coffin data.

```

28786 #4 { coffin ~ \__coffin_to_value:N #1 ~ corners }
28787   \l__coffin_corners_prop
28788 #4 { coffin ~ \__coffin_to_value:N #1 ~ poles }
28789   \l__coffin_poles_prop
28790 }

```

(End definition for \coffin_rotate:Nn, \coffin_grotate:Nn, and __coffin_rotate:NnNNN. These functions are documented on page 255.)

`__coffin_set_bounding:N` The bounding box corners for a coffin are easy enough to find: this is the same code as for the corners of the material itself, but using a dedicated property list.

```

28791 \cs_new_protected:Npn \__coffin_set_bounding:N #1
28792 {
28793   \prop_put:Nnx \l__coffin_bounding_prop { tl }
28794     { { 0pt } { \dim_eval:n { \box_ht:N #1 } } }
28795   \prop_put:Nnx \l__coffin_bounding_prop { tr }
28796     {
28797       { \dim_eval:n { \box_wd:N #1 } }

```

```

28798     { \dim_eval:n { \box_ht:N #1 } }
28799   }
28800   \dim_set:Nn \l__coffin_internal_dim { -\box_dp:N #1 }
28801   \prop_put:Nnx \l__coffin_bounding_prop { bl }
28802     { { Opt } { \dim_use:N \l__coffin_internal_dim } }
28803   \prop_put:Nnx \l__coffin_bounding_prop { br }
28804   {
28805     { \dim_eval:n { \box_wd:N #1 } }
28806     { \dim_use:N \l__coffin_internal_dim }
28807   }
28808 }

```

(End definition for `__coffin_set_bounding:N`.)

`__coffin_rotate_bounding:nnn` Rotating the position of the corner of the coffin is just a case of treating this as a vector
`__coffin_rotate_corner:Nnnn` from the reference point. The same treatment is used for the corners of the material itself
and the bounding box.

```

28809 \cs_new_protected:Npn \__coffin_rotate_bounding:nnn #1#2#3
28810 {
28811   \__coffin_rotate_vector:nnNN {#2} {#3} \l__coffin_x_dim \l__coffin_y_dim
28812   \prop_put:Nnx \l__coffin_bounding_prop {#1}
28813     { { \dim_use:N \l__coffin_x_dim } { \dim_use:N \l__coffin_y_dim } }
28814 }
28815 \cs_new_protected:Npn \__coffin_rotate_corner:Nnnn #1#2#3#4
28816 {
28817   \__coffin_rotate_vector:nnNN {#3} {#4} \l__coffin_x_dim \l__coffin_y_dim
28818   \prop_put:Nnx \l__coffin_corners_prop {#2}
28819     { { \dim_use:N \l__coffin_x_dim } { \dim_use:N \l__coffin_y_dim } }
28820 }

```

(End definition for `__coffin_rotate_bounding:nnn` and `__coffin_rotate_corner:Nnnn`.)

`__coffin_rotate_pole:Nnnnnn` Rotating a single pole simply means shifting the co-ordinate of the pole and its direction.
The rotation here is about the bottom-left corner of the coffin.

```

28821 \cs_new_protected:Npn \__coffin_rotate_pole:Nnnnnn #1#2#3#4#5#6
28822 {
28823   \__coffin_rotate_vector:nnNN {#3} {#4} \l__coffin_x_dim \l__coffin_y_dim
28824   \__coffin_rotate_vector:nnNN {#5} {#6}
28825     \l__coffin_x_prime_dim \l__coffin_y_prime_dim
28826   \prop_put:Nnx \l__coffin_poles_prop {#2}
28827   {
28828     { \dim_use:N \l__coffin_x_dim } { \dim_use:N \l__coffin_y_dim }
28829     { \dim_use:N \l__coffin_x_prime_dim }
28830     { \dim_use:N \l__coffin_y_prime_dim }
28831   }
28832 }

```

(End definition for `__coffin_rotate_pole:Nnnnnn`.)

`__coffin_rotate_vector:nnNN` A rotation function, which needs only an input vector (as dimensions) and an output
space. The values `\l__coffin_cos_fp` and `\l__coffin_sin_fp` should previously have
been set up correctly. Working this way means that the floating point work is kept to a
minimum: for any given rotation the sin and cosine values do no change, after all.

```

28833 \cs_new_protected:Npn \__coffin_rotate_vector:nnNN #1#2#3#4

```

```

28834 {
28835   \dim_set:Nn #3
28836   {
28837     \fp_to_dim:n
28838     {
28839       \dim_to_fp:n {#1} * \l__coffin_cos_fp
28840       - \dim_to_fp:n {#2} * \l__coffin_sin_fp
28841     }
28842   }
28843   \dim_set:Nn #4
28844   {
28845     \fp_to_dim:n
28846     {
28847       \dim_to_fp:n {#1} * \l__coffin_sin_fp
28848       + \dim_to_fp:n {#2} * \l__coffin_cos_fp
28849     }
28850   }
28851 }

```

(End definition for __coffin_rotate_vector:nnNN.)

_coffin_find_corner_maxima:N
_coffin_find_corner_maxima_aux:nn

The idea here is to find the extremities of the content of the coffin. This is done by looking for the smallest values for the bottom and left corners, and the largest values for the top and right corners. The values start at the maximum dimensions so that the case where all are positive or all are negative works out correctly.

```

28852 \cs_new_protected:Npn \__coffin_find_corner_maxima:N #1
28853 {
28854   \dim_set:Nn \l__coffin_top_corner_dim { -\c_max_dim }
28855   \dim_set:Nn \l__coffin_right_corner_dim { -\c_max_dim }
28856   \dim_set:Nn \l__coffin_bottom_corner_dim { \c_max_dim }
28857   \dim_set:Nn \l__coffin_left_corner_dim { \c_max_dim }
28858   \prop_map_inline:Nn \l__coffin_corners_prop
28859   { \__coffin_find_corner_maxima_aux:nn ##2 }
28860 }
28861 \cs_new_protected:Npn \__coffin_find_corner_maxima_aux:nn #1#2
28862 {
28863   \dim_set:Nn \l__coffin_left_corner_dim
28864   { \dim_min:nn { \l__coffin_left_corner_dim } {#1} }
28865   \dim_set:Nn \l__coffin_right_corner_dim
28866   { \dim_max:nn { \l__coffin_right_corner_dim } {#1} }
28867   \dim_set:Nn \l__coffin_bottom_corner_dim
28868   { \dim_min:nn { \l__coffin_bottom_corner_dim } {#2} }
28869   \dim_set:Nn \l__coffin_top_corner_dim
28870   { \dim_max:nn { \l__coffin_top_corner_dim } {#2} }
28871 }

```

(End definition for __coffin_find_corner_maxima:N and __coffin_find_corner_maxima_aux:nn.)

_coffin_find_bounding_shift:
_coffin_find_bounding_shift_aux:nn

The approach to finding the shift for the bounding box is similar to that for the corners. However, there is only one value needed here and a fixed input property list, so things are a bit clearer.

```

28872 \cs_new_protected:Npn \__coffin_find_bounding_shift:
28873 {
28874   \dim_set:Nn \l__coffin_bounding_shift_dim { \c_max_dim }

```

```

28875     \prop_map_inline:Nn \l__coffin_bounding_prop
28876     { \__coffin_find_bounding_shift_aux:nn ##2 }
28877   }
28878   \cs_new_protected:Npn \__coffin_find_bounding_shift_aux:nn #1#2
28879   {
28880     \dim_set:Nn \l__coffin_bounding_shift_dim
28881     { \dim_min:nn { \l__coffin_bounding_shift_dim } {#1} }
28882   }

```

(End definition for `__coffin_find_bounding_shift:` and `__coffin_find_bounding_shift_aux:nn`.)

`__coffin_shift_corner:Nnnn` `__coffin_shift_pole:Nnnnnn` Shifting the corners and poles of a coffin means subtracting the appropriate values from the x - and y -components. For the poles, this means that the direction vector is unchanged.

```

28883   \cs_new_protected:Npn \__coffin_shift_corner:Nnnn #1#2#3#4
28884   {
28885     \prop_put:Nnx \l__coffin_corners_prop {#2}
28886     {
28887       { \dim_eval:n { #3 - \l__coffin_left_corner_dim } }
28888       { \dim_eval:n { #4 - \l__coffin_bottom_corner_dim } }
28889     }
28890   }
28891   \cs_new_protected:Npn \__coffin_shift_pole:Nnnnnn #1#2#3#4#5#6
28892   {
28893     \prop_put:Nnx \l__coffin_poles_prop {#2}
28894     {
28895       { \dim_eval:n { #3 - \l__coffin_left_corner_dim } }
28896       { \dim_eval:n { #4 - \l__coffin_bottom_corner_dim } }
28897       {#5} {#6}
28898     }
28899   }

```

(End definition for `__coffin_shift_corner:Nnnn` and `__coffin_shift_pole:Nnnnnn`.)

`\l__coffin_scale_x_fp` `\l__coffin_scale_y_fp` Storage for the scaling factors in x and y , respectively.

```

28900   \fp_new:N \l__coffin_scale_x_fp
28901   \fp_new:N \l__coffin_scale_y_fp

```

(End definition for `\l__coffin_scale_x_fp` and `\l__coffin_scale_y_fp`.)

`\l__coffin_scaled_total_height_dim` `\l__coffin_scaled_width_dim` When scaling, the values given have to be turned into absolute values.

```

28902   \dim_new:N \l__coffin_scaled_total_height_dim
28903   \dim_new:N \l__coffin_scaled_width_dim

```

(End definition for `\l__coffin_scaled_total_height_dim` and `\l__coffin_scaled_width_dim`.)

`\coffin_resize:Nnn` `\coffin_resize:cnn` `\coffin_gresize:Nnn` `\coffin_gresize:cnn` Resizing a coffin begins by setting up the user-friendly names for the dimensions of the coffin box. The new sizes are then turned into scale factor. This is the same operation as takes place for the underlying box, but that operation is grouped and so the same calculation is done here.

```

\__coffin_resize:NnnNN
28904   \cs_new_protected:Npn \coffin_resize:Nnn #1#2#3
28905   {
28906     \__coffin_resize:NnnNN #1 {#2} {#3}
28907     \box_resize_to_wd_and_ht_plus_dp:Nnn

```

```

28908     \prop_set_eq:cN
28909   }
28910 \cs_generate_variant:Nn \coffin_resize:Nnn { c }
28911 \cs_new_protected:Npn \coffin_gresize:Nnn #1#2#3
28912 {
28913   \__coffin_resize:NnnNN #1 {#2} {#3}
28914   \box_gresize_to_wd_and_ht_plus_dp:Nnn
28915   \prop_gset_eq:cN
28916 }
28917 \cs_generate_variant:Nn \coffin_gresize:Nnn { c }
28918 \cs_new_protected:Npn \__coffin_resize:NnnNN #1#2#3#4#5
28919 {
28920   \fp_set:Nn \l__coffin_scale_x_fp
28921   { \dim_to_fp:n {#2} / \dim_to_fp:n { \coffin_wd:N #1 } }
28922   \fp_set:Nn \l__coffin_scale_y_fp
28923   {
28924     \dim_to_fp:n {#3}
28925     / \dim_to_fp:n { \coffin_ht:N #1 + \coffin_dp:N #1 }
28926   }
28927   #4 #1 {#2} {#3}
28928   \__coffin_resize_common:NnnN #1 {#2} {#3} #5
28929 }

```

(End definition for `\coffin_resize:Nnn`, `\coffin_gresize:Nnn`, and `__coffin_resize:NnnNN`. These functions are documented on page 255.)

`__coffin_resize_common:NnnN` The poles and corners of the coffin are scaled to the appropriate places before actually resizing the underlying box.

```

28930 \cs_new_protected:Npn \__coffin_resize_common:NnnN #1#2#3#4
28931 {
28932   \prop_set_eq:Nc \l__coffin_corners_prop
28933   { coffin ~ \__coffin_to_value:N #1 ~ corners }
28934   \prop_set_eq:Nc \l__coffin_poles_prop
28935   { coffin ~ \__coffin_to_value:N #1 ~ poles }
28936   \prop_map_inline:Nn \l__coffin_corners_prop
28937   { \__coffin_scale_corner:Nnnn #1 {##1} ##2 }
28938   \prop_map_inline:Nn \l__coffin_poles_prop
28939   { \__coffin_scale_pole:Nnnnnn #1 {##1} ##2 }

```

Negative x -scaling values place the poles in the wrong location: this is corrected here.

```

28940   \fp_compare:nNnT \l__coffin_scale_x_fp < \c_zero_fp
28941   {
28942     \prop_map_inline:Nn \l__coffin_corners_prop
28943     { \__coffin_x_shift_corner:Nnnn #1 {##1} ##2 }
28944     \prop_map_inline:Nn \l__coffin_poles_prop
28945     { \__coffin_x_shift_pole:Nnnnnn #1 {##1} ##2 }
28946   }
28947   #4 { coffin ~ \__coffin_to_value:N #1 ~ corners }
28948   \l__coffin_corners_prop
28949   #4 { coffin ~ \__coffin_to_value:N #1 ~ poles }
28950   \l__coffin_poles_prop
28951 }

```

(End definition for `__coffin_resize_common:NnnN`.)

`\coffin_scale:Nnn` For scaling, the opposite calculation is done to find the new dimensions for the coffin.
`\coffin_scale:cnn` Only the total height is needed, as this is the shift required for corners and poles. The
`\coffin_gscale:Nnn` scaling is done the T_EX way as this works properly with floating point values without
`\coffin_gscale:cnn` needing to use the fp module.
`\coffin_scale:NnnNN`

```

28952 \cs_new_protected:Npn \coffin_scale:Nnn #1#2#3
28953 { \__coffin_scale:NnnNN #1 {#2} {#3} \box_scale:Nnn \prop_set_eq:cN }
28954 \cs_generate_variant:Nn \coffin_scale:Nnn { c }
28955 \cs_new_protected:Npn \coffin_gscale:Nnn #1#2#3
28956 { \__coffin_scale:NnnNN #1 {#2} {#3} \box_gscale:Nnn \prop_gset_eq:cN }
28957 \cs_generate_variant:Nn \coffin_gscale:Nnn { c }
28958 \cs_new_protected:Npn \__coffin_scale:NnnNN #1#2#3#4#5
28959 {
28960   \fp_set:Nn \l__coffin_scale_x_fp {#2}
28961   \fp_set:Nn \l__coffin_scale_y_fp {#3}
28962   #4 #1 { \l__coffin_scale_x_fp } { \l__coffin_scale_y_fp }
28963   \dim_set:Nn \l__coffin_internal_dim
28964     { \coffin_ht:N #1 + \coffin_dp:N #1 }
28965   \dim_set:Nn \l__coffin_scaled_total_height_dim
28966     { \fp_abs:n { \l__coffin_scale_y_fp } \l__coffin_internal_dim }
28967   \dim_set:Nn \l__coffin_scaled_width_dim
28968     { -\fp_abs:n { \l__coffin_scale_x_fp } \coffin_wd:N #1 }
28969   \__coffin_resize_common:NnnN #1
28970     { \l__coffin_scaled_width_dim } { \l__coffin_scaled_total_height_dim }
28971   #5
28972 }

```

(End definition for `\coffin_scale:Nnn`, `\coffin_gscale:Nnn`, and `\coffin_scale:NnnNN`. These functions are documented on page 255.)

`__coffin_scale_vector:nnNN` This functions scales a vector from the origin using the pre-set scale factors in x and y . This is a much less complex operation than rotation, and as a result the code is a lot clearer.

```

28973 \cs_new_protected:Npn \__coffin_scale_vector:nnNN #1#2#3#4
28974 {
28975   \dim_set:Nn #3
28976     { \fp_to_dim:n { \dim_to_fp:n {#1} * \l__coffin_scale_x_fp } }
28977   \dim_set:Nn #4
28978     { \fp_to_dim:n { \dim_to_fp:n {#2} * \l__coffin_scale_y_fp } }
28979 }

```

(End definition for `__coffin_scale_vector:nnNN`.)

`__coffin_scale_corner:Nnnn` Scaling both corners and poles is a simple calculation using the preceding vector scaling.
`__coffin_scale_pole:Nnnnnn`

```

28980 \cs_new_protected:Npn \__coffin_scale_corner:Nnnn #1#2#3#4
28981 {
28982   \__coffin_scale_vector:nnNN {#3} {#4} \l__coffin_x_dim \l__coffin_y_dim
28983   \prop_put:Nnx \l__coffin_corners_prop {#2}
28984     { { \dim_use:N \l__coffin_x_dim } { \dim_use:N \l__coffin_y_dim } }
28985 }
28986 \cs_new_protected:Npn \__coffin_scale_pole:Nnnnnn #1#2#3#4#5#6
28987 {
28988   \__coffin_scale_vector:nnNN {#3} {#4} \l__coffin_x_dim \l__coffin_y_dim
28989   \prop_put:Nnx \l__coffin_poles_prop {#2}
28990   {

```

```

28991         { \dim_use:N \l__coffin_x_dim } { \dim_use:N \l__coffin_y_dim }
28992         {#5} {#6}
28993     }
28994 }

```

(End definition for `__coffin_scale_corner:Nnnn` and `__coffin_scale_pole:Nnnnnn`.)

`__coffin_x_shift_corner:Nnnn`
`__coffin_x_shift_pole:Nnnnnn`

These functions correct for the x displacement that takes place with a negative horizontal scaling.

```

28995 \cs_new_protected:Npn \__coffin_x_shift_corner:Nnnn #1#2#3#4
28996 {
28997     \prop_put:Nnx \l__coffin_corners_prop {#2}
28998     {
28999         { \dim_eval:n { #3 + \box_wd:N #1 } } {#4}
29000     }
29001 }
29002 \cs_new_protected:Npn \__coffin_x_shift_pole:Nnnnnn #1#2#3#4#5#6
29003 {
29004     \prop_put:Nnx \l__coffin_poles_prop {#2}
29005     {
29006         { \dim_eval:n { #3 + \box_wd:N #1 } } {#4}
29007         {#5} {#6}
29008     }
29009 }

```

(End definition for `__coffin_x_shift_corner:Nnnn` and `__coffin_x_shift_pole:Nnnnnn`.)

43.7 Aligning and typesetting of coffins

`\coffin_join:NnnNnnnn`
`\coffin_join:cnnNnnnn`
`\coffin_join:Nnncnnnn`
`\coffin_join:cnncnnnn`

This command joins two coffins, using a horizontal and vertical pole from each coffin and making an offset between the two. The result is stored as the as a third coffin, which has all of its handles reset to standard values. First, the more basic alignment function is used to get things started.

```

\coffin_gjoin:NnnNnnnn
\coffin_gjoin:cnnNnnnn
\coffin_gjoin:Nnncnnnn
\coffin_gjoin:cnncnnnn
\__coffin_join:NnnNnnnnN
29010 \cs_new_protected:Npn \coffin_join:NnnNnnnn #1#2#3#4#5#6#7#8
29011 {
29012     \__coffin_join:NnnNnnnnN #1 {#2} {#3} #4 {#5} {#6} {#7} {#8}
29013     \coffin_set_eq:NN
29014 }
29015 \cs_generate_variant:Nn \coffin_join:NnnNnnnn { c , Nnnc , cnnc }
29016 \cs_new_protected:Npn \coffin_gjoin:NnnNnnnn #1#2#3#4#5#6#7#8
29017 {
29018     \__coffin_join:NnnNnnnnN #1 {#2} {#3} #4 {#5} {#6} {#7} {#8}
29019     \coffin_gset_eq:NN
29020 }
29021 \cs_generate_variant:Nn \coffin_gjoin:NnnNnnnn { c , Nnnc , cnnc }
29022 \cs_new_protected:Npn \__coffin_join:NnnNnnnnN #1#2#3#4#5#6#7#8#9
29023 {
29024     \__coffin_align:NnnNnnnnN
29025     #1 {#2} {#3} #4 {#5} {#6} {#7} {#8} \l__coffin_aligned_coffin

```

Correct the placement of the reference point. If the x -offset is negative then the reference point of the second box is to the left of that of the first, which is corrected using a kern. On the right side the first box might stick out, which would show up if it is wider than the sum of the x -offset and the width of the second box. So a second kern may be needed.


```

29026 \hbox_set:Nn \l__coffin_aligned_coffin
29027 {
29028   \dim_compare:nNnT { \l__coffin_offset_x_dim } < \c_zero_dim
29029   { \tex_kern:D -\l__coffin_offset_x_dim }
29030   \hbox_unpack:N \l__coffin_aligned_coffin
29031   \dim_set:Nn \l__coffin_internal_dim
29032   { \l__coffin_offset_x_dim - \box_wd:N #1 + \box_wd:N #4 }
29033   \dim_compare:nNnT \l__coffin_internal_dim < \c_zero_dim
29034   { \tex_kern:D -\l__coffin_internal_dim }
29035 }

```

The coffin structure is reset, and the corners are cleared: only those from the two parent coffins are needed.

```

29036 \__coffin_reset_structure:N \l__coffin_aligned_coffin
29037 \prop_clear:c
29038 {
29039   coffin ~ \__coffin_to_value:N \l__coffin_aligned_coffin
29040   \c_space_tl corners
29041 }
29042 \__coffin_update_poles:N \l__coffin_aligned_coffin

```

The structures of the parent coffins are now transferred to the new coffin, which requires that the appropriate offsets are applied. That then depends on whether any shift was needed.

```

29043 \dim_compare:nNnTF \l__coffin_offset_x_dim < \c_zero_dim
29044 {
29045   \__coffin_offset_poles:Nnn #1 { -\l__coffin_offset_x_dim } { Opt }
29046   \__coffin_offset_poles:Nnn #4 { Opt } { \l__coffin_offset_y_dim }
29047   \__coffin_offset_corners:Nnn #1 { -\l__coffin_offset_x_dim } { Opt }
29048   \__coffin_offset_corners:Nnn #4 { Opt } { \l__coffin_offset_y_dim }
29049 }
29050 {
29051   \__coffin_offset_poles:Nnn #1 { Opt } { Opt }
29052   \__coffin_offset_poles:Nnn #4
29053   { \l__coffin_offset_x_dim } { \l__coffin_offset_y_dim }
29054   \__coffin_offset_corners:Nnn #1 { Opt } { Opt }
29055   \__coffin_offset_corners:Nnn #4
29056   { \l__coffin_offset_x_dim } { \l__coffin_offset_y_dim }
29057 }
29058 \__coffin_update_vertical_poles:NNN #1 #4 \l__coffin_aligned_coffin
29059 #9 #1 \l__coffin_aligned_coffin
29060 }

```

(End definition for `\coffin_join:NnnNnnnn`, `\coffin_gjoin:NnnNnnnn`, and `__coffin_join:NnnNnnnnN`. These functions are documented on page 255.)

```

\coffin_attach:NnnNnnnn
\coffin_attach:cnnNnnnn
\coffin_attach:Nnnncnnnn
\coffin_attach:cnncnnnn
\coffin_gattach:NnnNnnnn
\coffin_gattach:cnnNnnnn
\coffin_gattach:Nnnncnnnn
\coffin_gattach:cnncnnnn
\__coffin_attach:NnnNnnnnN
  \__coffin_attach_mark:NnnNnnnn

```

A more simple version of the above, as it simply uses the size of the first coffin for the new one. This means that the work here is rather simplified compared to the above code. The function used when marking a position is hear also as it is similar but without the structure updates.

```

29061 \cs_new_protected:Npn \coffin_attach:NnnNnnnn #1#2#3#4#5#6#7#8
29062 {
29063   \__coffin_attach:NnnNnnnnN #1 {#2} {#3} #4 {#5} {#6} {#7} {#8}
29064   \coffin_set_eq:NN
29065 }

```

```

29066 \cs_generate_variant:Nn \coffin_attach:NnnNnnnn { c , Nnnc , cnnc }
29067 \cs_new_protected:Npn \coffin_gattach:NnnNnnnn #1#2#3#4#5#6#7#8
29068 {
29069   \__coffin_attach:NnnNnnnnN #1 {#2} {#3} #4 {#5} {#6} {#7} {#8}
29070   \coffin_gset_eq:NN
29071 }
29072 \cs_generate_variant:Nn \coffin_gattach:NnnNnnnn { c , Nnnc , cnnc }
29073 \cs_new_protected:Npn \__coffin_attach:NnnNnnnnN #1#2#3#4#5#6#7#8#9
29074 {
29075   \__coffin_align:NnnNnnnnN
29076   #1 {#2} {#3} #4 {#5} {#6} {#7} {#8} \l__coffin_aligned_coffin
29077   \box_set_ht:Nn \l__coffin_aligned_coffin { \box_ht:N #1 }
29078   \box_set_dp:Nn \l__coffin_aligned_coffin { \box_dp:N #1 }
29079   \box_set_wd:Nn \l__coffin_aligned_coffin { \box_wd:N #1 }
29080   \__coffin_reset_structure:N \l__coffin_aligned_coffin
29081   \prop_set_eq:cc
29082   {
29083     coffin ~ \__coffin_to_value:N \l__coffin_aligned_coffin
29084     \c_space_tl corners
29085   }
29086   { coffin ~ \__coffin_to_value:N #1 ~ corners }
29087   \__coffin_update_poles:N \l__coffin_aligned_coffin
29088   \__coffin_offset_poles:Nnn #1 { Opt } { Opt }
29089   \__coffin_offset_poles:Nnn #4
29090   { \l__coffin_offset_x_dim } { \l__coffin_offset_y_dim }
29091   \__coffin_update_vertical_poles:NNN #1 #4 \l__coffin_aligned_coffin
29092   \coffin_set_eq:NN #1 \l__coffin_aligned_coffin
29093 }
29094 \cs_new_protected:Npn \__coffin_attach_mark:NnnNnnnn #1#2#3#4#5#6#7#8
29095 {
29096   \__coffin_align:NnnNnnnnN
29097   #1 {#2} {#3} #4 {#5} {#6} {#7} {#8} \l__coffin_aligned_coffin
29098   \box_set_ht:Nn \l__coffin_aligned_coffin { \box_ht:N #1 }
29099   \box_set_dp:Nn \l__coffin_aligned_coffin { \box_dp:N #1 }
29100   \box_set_wd:Nn \l__coffin_aligned_coffin { \box_wd:N #1 }
29101   \box_set_eq:NN #1 \l__coffin_aligned_coffin
29102 }

```

(End definition for \coffin_attach:NnnNnnnn and others. These functions are documented on page 255.)

`__coffin_align:NnnNnnnnN` The internal function aligns the two coffins into a third one, but performs no corrections on the resulting coffin poles. The process begins by finding the points of intersection for the poles for each of the input coffins. Those for the first coffin are worked out after those for the second coffin, as this allows the ‘primed’ storage area to be used for the second coffin. The ‘real’ box offsets are then calculated, before using these to re-box the input coffins. The default poles are then set up, but the final result depends on how the bounding box is being handled.

```

29103 \cs_new_protected:Npn \__coffin_align:NnnNnnnnN #1#2#3#4#5#6#7#8#9
29104 {
29105   \__coffin_calculate_intersection:Nnn #4 {#5} {#6}
29106   \dim_set:Nn \l__coffin_x_prime_dim { \l__coffin_x_dim }
29107   \dim_set:Nn \l__coffin_y_prime_dim { \l__coffin_y_dim }
29108   \__coffin_calculate_intersection:Nnn #1 {#2} {#3}

```

```

29109 \dim_set:Nn \l__coffin_offset_x_dim
29110 { \l__coffin_x_dim - \l__coffin_x_prime_dim + #7 }
29111 \dim_set:Nn \l__coffin_offset_y_dim
29112 { \l__coffin_y_dim - \l__coffin_y_prime_dim + #8 }
29113 \hbox_set:Nn \l__coffin_aligned_internal_coffin
29114 {
29115   \box_use:N #1
29116   \tex_kern:D -\box_wd:N #1
29117   \tex_kern:D \l__coffin_offset_x_dim
29118   \box_move_up:nn { \l__coffin_offset_y_dim } { \box_use:N #4 }
29119 }
29120 \coffin_set_eq:NN #9 \l__coffin_aligned_internal_coffin
29121 }

```

(End definition for _coffin_align:NnnNnnnnN.)

_coffin_offset_poles:Nnn
_coffin_offset_pole:Nnnnnnn

Transferring structures from one coffin to another requires that the positions are updated by the offset between the two coffins. This is done by mapping to the property list of the source coffins, moving as appropriate and saving to the new coffin data structures. The test for a - means that the structures from the parent coffins are uniquely labelled and do not depend on the order of alignment. The pay off for this is that - should not be used in coffin pole or handle names, and that multiple alignments do not result in a whole set of values.

```

29122 \cs_new_protected:Npn \_coffin_offset_poles:Nnn #1#2#3
29123 {
29124   \prop_map_inline:cn { coffin ~ \_coffin_to_value:N #1 ~ poles }
29125   { \_coffin_offset_pole:Nnnnnnn #1 {##1} ##2 {#2} {#3} }
29126 }
29127 \cs_new_protected:Npn \_coffin_offset_pole:Nnnnnnn #1#2#3#4#5#6#7#8
29128 {
29129   \dim_set:Nn \l__coffin_x_dim { #3 + #7 }
29130   \dim_set:Nn \l__coffin_y_dim { #4 + #8 }
29131   \tl_if_in:nnTF {#2} { - }
29132   { \tl_set:Nn \l__coffin_internal_tl { {#2} } }
29133   { \tl_set:Nn \l__coffin_internal_tl { { #1 - #2 } } }
29134   \exp_last_unbraced:NNo \_coffin_set_pole:Nnx \l__coffin_aligned_coffin
29135   { \l__coffin_internal_tl }
29136   {
29137     { \dim_use:N \l__coffin_x_dim } { \dim_use:N \l__coffin_y_dim }
29138     {#5} {#6}
29139   }
29140 }

```

(End definition for _coffin_offset_poles:Nnn and _coffin_offset_pole:Nnnnnnn.)

_coffin_offset_corners:Nnn
_coffin_offset_corner:Nnnnn

Saving the offset corners of a coffin is very similar, except that there is no need to worry about naming: every corner can be saved here as order is unimportant.

```

29141 \cs_new_protected:Npn \_coffin_offset_corners:Nnn #1#2#3
29142 {
29143   \prop_map_inline:cn { coffin ~ \_coffin_to_value:N #1 ~ corners }
29144   { \_coffin_offset_corner:Nnnnn #1 {##1} ##2 {#2} {#3} }
29145 }
29146 \cs_new_protected:Npn \_coffin_offset_corner:Nnnnn #1#2#3#4#5#6
29147 {

```

```

29148 \prop_put:cnx
29149 {
29150   coffin ~ \__coffin_to_value:N \l__coffin_aligned_coffin
29151   \c_space_tl corners
29152 }
29153 { #1 - #2 }
29154 {
29155   { \dim_eval:n { #3 + #5 } }
29156   { \dim_eval:n { #4 + #6 } }
29157 }
29158 }

```

(End definition for __coffin_offset_corners:Nnn and __coffin_offset_corner:Nnnnn.)

```

\__coffin_update_vertical_poles:NNN
\__coffin_update_T:nnnnnnnnN
\__coffin_update_B:nnnnnnnnN

```

The T and B poles need to be recalculated after alignment. These functions find the larger absolute value for the poles, but this is of course only logical when the poles are horizontal.

```

29159 \cs_new_protected:Npn \__coffin_update_vertical_poles:NNN #1#2#3
29160 {
29161   \__coffin_get_pole:NnN #3 { #1 -T } \l__coffin_pole_a_tl
29162   \__coffin_get_pole:NnN #3 { #2 -T } \l__coffin_pole_b_tl
29163   \exp_last_two_unbraced:Noo \__coffin_update_T:nnnnnnnnN
29164   \l__coffin_pole_a_tl \l__coffin_pole_b_tl #3
29165   \__coffin_get_pole:NnN #3 { #1 -B } \l__coffin_pole_a_tl
29166   \__coffin_get_pole:NnN #3 { #2 -B } \l__coffin_pole_b_tl
29167   \exp_last_two_unbraced:Noo \__coffin_update_B:nnnnnnnnN
29168   \l__coffin_pole_a_tl \l__coffin_pole_b_tl #3
29169 }
29170 \cs_new_protected:Npn \__coffin_update_T:nnnnnnnnN #1#2#3#4#5#6#7#8#9
29171 {
29172   \dim_compare:nNnTF {#2} < {#6}
29173   {
29174     \__coffin_set_pole:Nnx #9 { T }
29175     { { Opt } {#6} { 1000pt } { Opt } }
29176   }
29177   {
29178     \__coffin_set_pole:Nnx #9 { T }
29179     { { Opt } {#2} { 1000pt } { Opt } }
29180   }
29181 }
29182 \cs_new_protected:Npn \__coffin_update_B:nnnnnnnnN #1#2#3#4#5#6#7#8#9
29183 {
29184   \dim_compare:nNnTF {#2} < {#6}
29185   {
29186     \__coffin_set_pole:Nnx #9 { B }
29187     { { Opt } {#2} { 1000pt } { Opt } }
29188   }
29189   {
29190     \__coffin_set_pole:Nnx #9 { B }
29191     { { Opt } {#6} { 1000pt } { Opt } }
29192   }
29193 }

```

(End definition for __coffin_update_vertical_poles:NNN, __coffin_update_T:nnnnnnnnN, and __coffin_update_B:nnnnnnnnN.)

`\c__coffin_empty_coffin` An empty-but-horizontal coffin.

```

29194 \coffin_new:N \c__coffin_empty_coffin
29195 \tex_setbox:D \c__coffin_empty_coffin = \tex_hbox:D { }

(End definition for \c__coffin_empty_coffin.)

```

`\coffin_typeset:Nnnnn` Typesetting a coffin means aligning it with the current position, which is done using a coffin with no content at all. As well as aligning to the empty coffin, there is also a need to leave vertical mode, if necessary.

`\coffin_typeset:cnnnn`

```

29196 \cs_new_protected:Npn \coffin_typeset:Nnnnn #1#2#3#4#5
29197 {
29198   \mode_leave_vertical:
29199   \__coffin_align:NnnNnnnnN \c__coffin_empty_coffin { H } { 1 }
29200   #1 {#2} {#3} {#4} {#5} \l__coffin_aligned_coffin
29201   \box_use_drop:N \l__coffin_aligned_coffin
29202 }
29203 \cs_generate_variant:Nn \coffin_typeset:Nnnnn { c }

(End definition for \coffin_typeset:Nnnnn. This function is documented on page 256.)

```

43.8 Coffin diagnostics

`\l__coffin_display_coffin` Used for printing coffins with data structures attached.

```

\l__coffin_display_coord_coffin 29204 \coffin_new:N \l__coffin_display_coffin
\l__coffin_display_pole_coffin 29205 \coffin_new:N \l__coffin_display_coord_coffin
29206 \coffin_new:N \l__coffin_display_pole_coffin

```

(End definition for `\l__coffin_display_coffin`, `\l__coffin_display_coord_coffin`, and `\l__coffin_display_pole_coffin`.)

`\l__coffin_display_handles_prop` This property list is used to print coffin handles at suitable positions. The offsets are expressed as multiples of the basic offset value, which therefore acts as a scale-factor.

```

29207 \prop_new:N \l__coffin_display_handles_prop
29208 \prop_put:Nnn \l__coffin_display_handles_prop { tl }
29209 { { b } { r } { -1 } { 1 } }
29210 \prop_put:Nnn \l__coffin_display_handles_prop { thc }
29211 { { b } { hc } { 0 } { 1 } }
29212 \prop_put:Nnn \l__coffin_display_handles_prop { tr }
29213 { { b } { l } { 1 } { 1 } }
29214 \prop_put:Nnn \l__coffin_display_handles_prop { vcl }
29215 { { vc } { r } { -1 } { 0 } }
29216 \prop_put:Nnn \l__coffin_display_handles_prop { vhc }
29217 { { vc } { hc } { 0 } { 0 } }
29218 \prop_put:Nnn \l__coffin_display_handles_prop { vcr }
29219 { { vc } { l } { 1 } { 0 } }
29220 \prop_put:Nnn \l__coffin_display_handles_prop { bl }
29221 { { t } { r } { -1 } { -1 } }
29222 \prop_put:Nnn \l__coffin_display_handles_prop { bhc }
29223 { { t } { hc } { 0 } { -1 } }
29224 \prop_put:Nnn \l__coffin_display_handles_prop { br }
29225 { { t } { l } { 1 } { -1 } }
29226 \prop_put:Nnn \l__coffin_display_handles_prop { Tl }
29227 { { t } { r } { -1 } { -1 } }
29228 \prop_put:Nnn \l__coffin_display_handles_prop { Thc }

```

```

29229 { { t } { hc } { 0 } { -1 } }
29230 \prop_put:Nnn \l__coffin_display_handles_prop { Tr }
29231 { { t } { l } { 1 } { -1 } }
29232 \prop_put:Nnn \l__coffin_display_handles_prop { Hl }
29233 { { vc } { r } { -1 } { 1 } }
29234 \prop_put:Nnn \l__coffin_display_handles_prop { Hhc }
29235 { { vc } { hc } { 0 } { 1 } }
29236 \prop_put:Nnn \l__coffin_display_handles_prop { Hr }
29237 { { vc } { l } { 1 } { 1 } }
29238 \prop_put:Nnn \l__coffin_display_handles_prop { Bl }
29239 { { b } { r } { -1 } { -1 } }
29240 \prop_put:Nnn \l__coffin_display_handles_prop { Bhc }
29241 { { b } { hc } { 0 } { -1 } }
29242 \prop_put:Nnn \l__coffin_display_handles_prop { Br }
29243 { { b } { l } { 1 } { -1 } }

```

(End definition for \l__coffin_display_handles_prop.)

`\l__coffin_display_offset_dim` The standard offset for the label from the handle position when displaying handles.

```

29244 \dim_new:N \l__coffin_display_offset_dim
29245 \dim_set:Nn \l__coffin_display_offset_dim { 2pt }

```

(End definition for \l__coffin_display_offset_dim.)

`\l__coffin_display_x_dim` As the intersections of poles have to be calculated to find which ones to print, there is
`\l__coffin_display_y_dim` a need to avoid repetition. This is done by saving the intersection into two dedicated
values.

```

29246 \dim_new:N \l__coffin_display_x_dim
29247 \dim_new:N \l__coffin_display_y_dim

```

(End definition for \l__coffin_display_x_dim and \l__coffin_display_y_dim.)

`\l__coffin_display_poles_prop` A property list for printing poles: various things need to be deleted from this to get a
“nice” output.

```

29248 \prop_new:N \l__coffin_display_poles_prop

```

(End definition for \l__coffin_display_poles_prop.)

`\l__coffin_display_font_tl` Stores the settings used to print coffin data: this keeps things flexible.

```

29249 \tl_new:N \l__coffin_display_font_tl
29250 \bool_lazy_and:nnT
29251 { \cs_if_exist_p:N \fmtname }
29252 { \str_if_eq_p:Vn \fmtname { LaTeX2e } }
29253 {
29254 \tl_set:Nn \l__coffin_display_font_tl
29255 { \sffamily \tiny }
29256 }

```

(End definition for \l__coffin_display_font_tl.)

`__coffin_color:n` Calls `\color`, and otherwise does nothing if `\color` is not defined. As this is relatively
rarely used, we have it self-define rather than delay using a hook.

```

29257 \cs_new_protected:Npn \__coffin_color:n #1 {#1}
29258 \bool_lazy_and:nnT
29259 { \cs_if_exist_p:N \fmtname }

```

```

29260 { \str_if_eq_p:Vn \fmtname { LaTeX2e } }
29261 {
29262   \cs_gset_protected:Npn \__coffin_color:n
29263   {
29264     \cs_gset_protected:Npx \__coffin_color:n
29265     {
29266       \cs_if_exist:NTF \color_select:n
29267       { \color_select:n }
29268       {
29269         \cs_if_exist:NTF \color
29270         { \exp_not:N \color }
29271         { \exp_not:N \use_none:n }
29272       }
29273     }
29274   \__coffin_color:n
29275 }
29276 }

```

(End definition for __coffin_color:n.)

__coffin_rule:nn Abstract out creation of rules here until there is a higher-level interface.

```

29277 \cs_new_protected:Npn \__coffin_rule:nn #1#2
29278 {
29279   \mode_leave_vertical:
29280   \hbox:n { \tex_vrule:D width #1 height #2 \scan_stop: }
29281 }

```

(End definition for __coffin_rule:nn.)

\coffin_mark_handle:Nnnn Marking a single handle is relatively easy. The standard attachment function is used, meaning that there are two calculations for the location. However, this is likely to be okay given the load expected. Contrast with the more optimised version for showing all handles which comes next.

\coffin_mark_handle:cnnn

__coffin_mark_handle_aux:nnnnNnn

```

29282 \cs_new_protected:Npn \coffin_mark_handle:Nnnn #1#2#3#4
29283 {
29284   \hcoffin_set:Nn \l__coffin_display_pole_coffin
29285   {
29286     \__coffin_color:n {#4}
29287     \__coffin_rule:nn { 1pt } { 1pt }
29288   }
29289   \__coffin_attach_mark:NnnNnnnn #1 {#2} {#3}
29290   \l__coffin_display_pole_coffin { hc } { vc } { Opt } { Opt }
29291   \hcoffin_set:Nn \l__coffin_display_coord_coffin
29292   {
29293     \__coffin_color:n {#4}
29294     \l__coffin_display_font_tl
29295     ( \tl_to_str:n { #2 , #3 } )
29296   }
29297   \prop_get:NnN \l__coffin_display_handles_prop
29298   { #2 #3 } \l__coffin_internal_tl
29299   \quark_if_no_value:NTF \l__coffin_internal_tl
29300   {
29301     \prop_get:NnN \l__coffin_display_handles_prop
29302     { #3 #2 } \l__coffin_internal_tl

```

```

29303     \quark_if_no_value:NTF \l__coffin_internal_tl
29304     {
29305         \__coffin_attach_mark:NnnNnnnn #1 {#2} {#3}
29306         \l__coffin_display_coord_coffin { l } { vc }
29307         { 1pt } { 0pt }
29308     }
29309     {
29310         \exp_last_unbraced:No \__coffin_mark_handle_aux:nnnnNnn
29311         \l__coffin_internal_tl #1 {#2} {#3}
29312     }
29313 }
29314 {
29315     \exp_last_unbraced:No \__coffin_mark_handle_aux:nnnnNnn
29316     \l__coffin_internal_tl #1 {#2} {#3}
29317 }
29318 }
29319 \cs_new_protected:Npn \__coffin_mark_handle_aux:nnnnNnn #1#2#3#4#5#6#7
29320 {
29321     \__coffin_attach_mark:NnnNnnnn #5 {#6} {#7}
29322     \l__coffin_display_coord_coffin {#1} {#2}
29323     { #3 \l__coffin_display_offset_dim }
29324     { #4 \l__coffin_display_offset_dim }
29325 }
29326 \cs_generate_variant:Nn \coffin_mark_handle:Nnnn { c }

```

(End definition for `\coffin_mark_handle:Nnnn` and `__coffin_mark_handle_aux:nnnnNnn`. This function is documented on page 256.)

\coffin_display_handles:Nn
\coffin_display_handles:cn
`__coffin_display_handles_aux:nnnnnn`
`__coffin_display_handles_aux:nnnn`
`__coffin_display_attach:Nnnnn`

Printing the poles starts by removing any duplicates, for which the H poles is used as the definitive version for the baseline and bottom. Two loops are then used to find the combinations of handles for all of these poles. This is done such that poles are removed during the loops to avoid duplication.

```

29327 \cs_new_protected:Npn \coffin_display_handles:Nn #1#2
29328 {
29329     \hcoffin_set:Nn \l__coffin_display_pole_coffin
29330     {
29331         \__coffin_color:n {#2}
29332         \__coffin_rule:nn { 1pt } { 1pt }
29333     }
29334     \prop_set_eq:Nc \l__coffin_display_poles_prop
29335     { coffin ~ \__coffin_to_value:N #1 ~ poles }
29336     \__coffin_get_pole:NnN #1 { H } \l__coffin_pole_a_tl
29337     \__coffin_get_pole:NnN #1 { T } \l__coffin_pole_b_tl
29338     \tl_if_eq:NNT \l__coffin_pole_a_tl \l__coffin_pole_b_tl
29339     { \prop_remove:Nn \l__coffin_display_poles_prop { T } }
29340     \__coffin_get_pole:NnN #1 { B } \l__coffin_pole_b_tl
29341     \tl_if_eq:NNT \l__coffin_pole_a_tl \l__coffin_pole_b_tl
29342     { \prop_remove:Nn \l__coffin_display_poles_prop { B } }
29343     \coffin_set_eq:NN \l__coffin_display_coffin #1
29344     \prop_map_inline:Nn \l__coffin_display_poles_prop
29345     {
29346         \prop_remove:Nn \l__coffin_display_poles_prop {##1}
29347         \__coffin_display_handles_aux:nnnnnn {##1} ##2 {#2}
29348     }

```



```

29349     \box_use_drop:N \l__coffin_display_coffin
29350 }

```

For each pole there is a check for an intersection, which here does not give an error if none is found. The successful values are stored and used to align the pole coffin with the main coffin for output. The positions are recovered from the preset list if available.

```

29351 \cs_new_protected:Npn \__coffin_display_handles_aux:nnnnnn #1#2#3#4#5#6
29352 {
29353     \prop_map_inline:Nn \l__coffin_display_poles_prop
29354     {
29355         \bool_set_false:N \l__coffin_error_bool
29356         \__coffin_calculate_intersection:nnnnnnnn {#2} {#3} {#4} {#5} ##2
29357         \bool_if:NF \l__coffin_error_bool
29358         {
29359             \dim_set:Nn \l__coffin_display_x_dim { \l__coffin_x_dim }
29360             \dim_set:Nn \l__coffin_display_y_dim { \l__coffin_y_dim }
29361             \__coffin_display_attach:Nnnnn
29362             \l__coffin_display_pole_coffin { hc } { vc }
29363             { Opt } { Opt }
29364             \hcoffin_set:Nn \l__coffin_display_coord_coffin
29365             {
29366                 \__coffin_color:n {#6}
29367                 \l__coffin_display_font_tl
29368                 ( \tl_to_str:n { #1 , ##1 } )
29369             }
29370             \prop_get:NnN \l__coffin_display_handles_prop
29371             { #1 ##1 } \l__coffin_internal_tl
29372             \quark_if_no_value:NTF \l__coffin_internal_tl
29373             {
29374                 \prop_get:NnN \l__coffin_display_handles_prop
29375                 { ##1 #1 } \l__coffin_internal_tl
29376                 \quark_if_no_value:NTF \l__coffin_internal_tl
29377                 {
29378                     \__coffin_display_attach:Nnnnn
29379                     \l__coffin_display_coord_coffin { l } { vc }
29380                     { 1pt } { Opt }
29381                 }
29382                 {
29383                     \exp_last_unbraced:No
29384                     \__coffin_display_handles_aux:nnnn
29385                     \l__coffin_internal_tl
29386                 }
29387             }
29388             {
29389                 \exp_last_unbraced:No \__coffin_display_handles_aux:nnnn
29390                 \l__coffin_internal_tl
29391             }
29392         }
29393     }
29394 }
29395 \cs_new_protected:Npn \__coffin_display_handles_aux:nnnn #1#2#3#4
29396 {
29397     \__coffin_display_attach:Nnnnn
29398     \l__coffin_display_coord_coffin {#1} {#2}

```

```

29399     { #3 \l__coffin_display_offset_dim }
29400     { #4 \l__coffin_display_offset_dim }
29401   }
29402 \cs_generate_variant:Nn \coffin_display_handles:Nn { c }

```

This is a dedicated version of `\coffin_attach:NnnNnnnn` with a hard-wired first coffin. As the intersection is already known and stored for the display coffin the code simply uses it directly, with no calculation.

```

29403 \cs_new_protected:Npn \__coffin_display_attach:Nnnnn #1#2#3#4#5
29404 {
29405   \__coffin_calculate_intersection:Nnn #1 {#2} {#3}
29406   \dim_set:Nn \l__coffin_x_prime_dim { \l__coffin_x_dim }
29407   \dim_set:Nn \l__coffin_y_prime_dim { \l__coffin_y_dim }
29408   \dim_set:Nn \l__coffin_offset_x_dim
29409     { \l__coffin_display_x_dim - \l__coffin_x_prime_dim + #4 }
29410   \dim_set:Nn \l__coffin_offset_y_dim
29411     { \l__coffin_display_y_dim - \l__coffin_y_prime_dim + #5 }
29412   \hbox_set:Nn \l__coffin_aligned_coffin
29413     {
29414       \box_use:N \l__coffin_display_coffin
29415       \tex_kern:D -\box_wd:N \l__coffin_display_coffin
29416       \tex_kern:D \l__coffin_offset_x_dim
29417       \box_move_up:nn { \l__coffin_offset_y_dim } { \box_use:N #1 }
29418     }
29419   \box_set_ht:Nn \l__coffin_aligned_coffin
29420     { \box_ht:N \l__coffin_display_coffin }
29421   \box_set_dp:Nn \l__coffin_aligned_coffin
29422     { \box_dp:N \l__coffin_display_coffin }
29423   \box_set_wd:Nn \l__coffin_aligned_coffin
29424     { \box_wd:N \l__coffin_display_coffin }
29425   \box_set_eq:NN \l__coffin_display_coffin \l__coffin_aligned_coffin
29426 }

```

(End definition for `\coffin_display_handles:Nn` and others. This function is documented on page 256.)

```

\coffin_show_structure:N
\coffin_show_structure:c
\coffin_log_structure:N
\coffin_log_structure:c
\__coffin_show_structure:NN

```

For showing the various internal structures attached to a coffin in a way that keeps things relatively readable. If there is no apparent structure then the code complains.

```

29427 \cs_new_protected:Npn \coffin_show_structure:N
29428   { \__coffin_show_structure:NN \msg_show:nnxxxx }
29429 \cs_generate_variant:Nn \coffin_show_structure:N { c }
29430 \cs_new_protected:Npn \coffin_log_structure:N
29431   { \__coffin_show_structure:NN \msg_log:nnxxxx }
29432 \cs_generate_variant:Nn \coffin_log_structure:N { c }
29433 \cs_new_protected:Npn \__coffin_show_structure:NN #1#2
29434   {
29435     \__coffin_if_exist:NT #2
29436     {
29437       #1 { LaTeX / kernel } { show-coffin }
29438       { \token_to_str:N #2 }
29439       {
29440         \iow_newline: >~ ht ~~~ \dim_eval:n { \coffin_ht:N #2 }
29441         \iow_newline: >~ dp ~~~ \dim_eval:n { \coffin_dp:N #2 }
29442         \iow_newline: >~ wd ~~~ \dim_eval:n { \coffin_wd:N #2 }
29443       }

```

```

29444         {
29445             \prop_map_function:cN
29446             { coffin ~ \__coffin_to_value:N #2 ~ poles }
29447             \msg_show_item_unbraced:nn
29448         }
29449     { }
29450 }
29451 }

```

(End definition for `\coffin_show_structure:N`, `\coffin_log_structure:N`, and `__coffin_show_structure:NN`. These functions are documented on page 256.)

43.9 Messages

```

29452 \__kernel_msg_new:nnnn { kernel } { no-pole-intersection }
29453 { No~intersection~between~coffin~poles. }
29454 {
29455     LaTeX~was~asked~to~find~the~intersection~between~two~poles,~
29456     but~they~do~not~have~a~unique~meeting~point:~
29457     the~value~(Opt,~Opt)~will~be~used.
29458 }
29459 \__kernel_msg_new:nnnn { kernel } { unknown-coffin }
29460 { Unknown~coffin~'#1'. }
29461 { The~coffin~'#1'~was~never~defined. }
29462 \__kernel_msg_new:nnnn { kernel } { unknown-coffin-pole }
29463 { Pole~'#1'~unknown~for~coffin~'#2'. }
29464 {
29465     LaTeX~was~asked~to~find~a~typesetting~pole~for~a~coffin,~
29466     but~either~the~coffin~does~not~exist~or~the~pole~name~is~wrong.
29467 }
29468 \__kernel_msg_new:nnn { kernel } { show-coffin }
29469 {
29470     Size~of~coffin~#1 : #2 \\
29471     Poles~of~coffin~#1 : #3 .
29472 }
29473 </package>

```

44 l3color-base Implementation

```

29474 <*package>
29475 <@@=color>

```

`\l__color_current_tl` The color currently active for foreground (text, *etc.*) material. This is stored in the form of a color model followed by one or more values. There are four pre-defined models, three of which take numerical values in the range [0, 1]:

- `gray` `<gray>` Grayscale color with the `<gray>` value running from 0 (fully black) to 1 (fully white)
- `cmymk` `<cyan>` `<magenta>` `<yellow>` `<black>`
- `rgb` `<red>` `<green>` `<blue>`

Notice that the value are separated by spaces. There is a fourth pre-defined model using a string value and a numerical one:

- **spot** $\langle name \rangle$ $\langle tint \rangle$ A pre-defined spot color, where the $\langle name \rangle$ should be a pre-defined string color name and the $\langle tint \rangle$ should be in the range $[0, 1]$.

Additional models may be created to allow mixing of spot colors. The number of data entries these require will depend on the number of colors to be mixed.

TeXhackers note: The content of `\l__color_current_tl` comprises two brace groups, the first containing the color model and the second containing the value(s) applicable in that model.

(End definition for \l__color_current_tl.)

\color_group_begin: Grouping for color is the same as using the basic `\group_begin:` and `\group_end:` functions. However, for semantic reasons, they are renamed here.

```
29476 \cs_new_eq:NN \color_group_begin: \group_begin:
29477 \cs_new_eq:NN \color_group_end: \group_end:
```

(End definition for \color_group_begin: and \color_group_end:.. These functions are documented on page 258.)

\color_ensure_current: A driver-independent wrapper for setting the foreground color to the current color “now”.

```
29478 \cs_new_protected:Npn \color_ensure_current:
29479 {
29480   \__color_backend_pickup:N \l__color_current_tl
29481   \__color_select:N \l__color_current_tl
29482 }
```

(End definition for \color_ensure_current:.. This function is documented on page 258.)

\s__color_stop Internal scan marks.

```
29483 \scan_new:N \s__color_stop
```

(End definition for \s__color_stop.)

__color_select:N Take an internal color specification and pass it to the driver. This code is needed to ensure the current color but will also be used by the higher-level experimental material.

__color_select:nn

```
29484 \cs_new_protected:Npn \__color_select:N #1
29485 { \exp_after:wN \__color_select:nn #1 }
29486 \cs_new_protected:Npn \__color_select:nn #1#2
29487 { \use:c { \__color_backend_select_ #1 :n } {#2} }
```

(End definition for __color_select:N and __color_select:nn.)

\l__color_current_tl The current color, with the model and

```
29488 \tl_new:N \l__color_current_tl
29489 \tl_set:Nn \l__color_current_tl { { gray } { 0 } }
```

(End definition for \l__color_current_tl.)

```
29490 \endpackage
```

45 l3luatex implementation

29491 `*package`

45.1 Breaking out to Lua

29492 `*tex`

29493 `\@@=lua`

29494 `\cs_new_eq:NN __lua_escape:n \tex_luaescapestring:D`
 29495 `\cs_new_eq:NN __lua_now:n \tex_directlua:D`
 29496 `\cs_new_eq:NN __lua_shipout:n \tex_latelua:D`

Copies of primitives.

(End definition for `__lua_escape:n`, `__lua_now:n`, and `__lua_shipout:n`.)

These functions are set up in `l3str` for bootstrapping: we want to replace them with a “proper” version at this stage, so clean up.

29497 `\cs_undefine:N \lua_escape:e`
 29498 `\cs_undefine:N \lua_now:e`

`\lua_now:n` Wrappers around the primitives. As with engines other than LuaTeX these have to be macros, we give them the same status in all cases. When LuaTeX is not in use, simply give an error message/
`\lua_now:e`
`\lua_shipout_e:n`
`\lua_shipout:n`
`\lua_escape:n`
`\lua_escape:e`

29499 `\cs_new:Npn \lua_now:e #1 { __lua_now:n {#1} }`
 29500 `\cs_new:Npn \lua_now:n #1 { \lua_now:e { \exp_not:n {#1} } }`
 29501 `\cs_new_protected:Npn \lua_shipout_e:n #1 { __lua_shipout:n {#1} }`
 29502 `\cs_new_protected:Npn \lua_shipout:n #1`
 29503 `{ \lua_shipout_e:n { \exp_not:n {#1} } }`
 29504 `\cs_new:Npn \lua_escape:e #1 { __lua_escape:n {#1} }`
 29505 `\cs_new:Npn \lua_escape:n #1 { \lua_escape:e { \exp_not:n {#1} } }`
 29506 `\sys_if_engine luatex:F`
 29507 `{`
 29508 `\clist_map_inline:nn`
 29509 `{`
 29510 `\lua_escape:n , \lua_escape:e ,`
 29511 `\lua_now:n , \lua_now:e`
 29512 `}`
 29513 `{`
 29514 `\cs_set:Npn #1 ##1`
 29515 `{`
 29516 `__kernel_msg_expandable_error:nnn`
 29517 `{ kernel } { luatex-required } { #1 }`
 29518 `}`
 29519 `}`
 29520 `\clist_map_inline:nn`
 29521 `{ \lua_shipout_e:n , \lua_shipout:n }`
 29522 `{`
 29523 `\cs_set_protected:Npn #1 ##1`
 29524 `{`
 29525 `__kernel_msg_error:nnn`
 29526 `{ kernel } { luatex-required } { #1 }`
 29527 `}`
 29528 `}`
 29529 `}`

(End definition for `\lua_now:n` and others. These functions are documented on page 259.)

45.2 Messages

```
29530 \__kernel_msg_new:nnnn { kernel } { luatex-required }
29531 { LuaTeX~engine~not~in~use!~Ignoring~#1. }
29532 {
29533     The~feature~you~are~using~is~only~available~
29534     with~the~LuaTeX~engine.~LaTeX3~ignored~'~#1'~.
29535 }
29536 </tex>
```

45.3 Lua functions for internal use

```
29537 <*lua>
```

Most of the emulation of pdfTEX here is based heavily on Heiko Oberdiek's pdfTeX-cmds package.

Create a table for the kernel's own use.

```
13kernel
ltx.utils
29538 l3kernel = l3kernel or { }
29539 local l3kernel = l3kernel
29540 ltx = ltx or {utils={}}
29541 ltx.utils = ltx.utils or { }
29542 local ltxutils = ltx.utils
```

(End definition for l3kernel and ltx.utils. These functions are documented on page 260.)

Local copies of global tables.

```
29543 local io      = io
29544 local kpse    = kpse
29545 local lfs     = lfs
29546 local math    = math
29547 local md5     = md5
29548 local os      = os
29549 local string  = string
29550 local tex     = tex
29551 local texio   = texio
29552 local tonumber = tonumber
```

Local copies of standard functions.

```
29553 local abs      = math.abs
29554 local byte     = string.byte
29555 local floor    = math.floor
29556 local format   = string.format
29557 local gsub     = string.gsub
29558 local lfs_attr = lfs.attributes
29559 local open     = io.open
29560 local os_date  = os.date
29561 local setcatcode = tex.setcatcode
29562 local sprint   = tex.sprint
29563 local cprint   = tex.cprint
29564 local write    = tex.write
29565 local write_nl = texio.write_nl
29566 local utf8_char = utf8.char
29567
29568 local scan_int    = token.scan_int or token.scan_integer
29569 local scan_string = token.scan_string
29570 local scan_keyword = token.scan_keyword
```

```

29571 local put_next      = token.put_next
29572
29573 local true_tok       = token.create'prg_return_true:'
29574 local false_tok      = token.create'prg_return_false:'
29575
29576 local function deprecated(table, name, func)
29577     table[name] = function(...)
29578         write_nl(format("Calling deprecated Lua function %s", name))
29579         table[name] = func
29580         return func(...)
29581     end
29582 end
29583 % \end{macrocode}
29584 % Deal with Con\TeX{}t: doesn't use |kpse| library.
29585 % \begin{macrocode}
29586 local kpse_find = (resolvers and resolvers.findfile) or kpse.find_file

```

escapehex An internal auxiliary to convert a string to the matching hex escape. This works on a byte basis: extension to handled UTF-8 input is covered in `pdftexcmds` but is not currently required here.

```

29587 local function escapehex(str)
29588     return (gsub(str, ".",
29589         function (ch) return format("%02X", byte(ch)) end))
29590 end

```

(End definition for `escapehex`.)

l3kernel.charcat Creating arbitrary chars using `tex.cprint`. The alternative approach using `token.put_next(token.cre` would be about 10% slower.

```

29591 deprecated(l3kernel, 'charcat', function(charcode, catcode)
29592     cprint(catcode, utf8_char(charcode))
29593 end)

```

(End definition for `l3kernel.charcat`. This function is documented on page 260.)

l3kernel.elapsedtime Simple timing set up: give the result from the system clock in scaled seconds.

l3kernel.resettimer

```

29594 local os_clock      = os.clock
29595 local base_clock_time = 0
29596 local function elapsedtime()
29597     local val = (os_clock() - base_clock_time) * 65536 + 0.5
29598     if val > 2147483647 then
29599         val = 2147483647
29600     end
29601     write(format("%d", floor(val)))
29602 end
29603 l3kernel.elapsedtime = elapsedtime
29604 local function resettimer()
29605     base_clock_time = os_clock()
29606 end
29607 l3kernel.resettimer = resettimer

```

(End definition for `l3kernel.elapsedtime` and `l3kernel.resettimer`. These functions are documented on page 260.)

`ltx.utils.filedump` Similar comments here to the next function: read the file in binary mode to avoid any
`l3kernel.filedump` line-end weirdness.

```

29608 local function filedump(name,offset,length)
29609   local file = kpse_find(name,"tex",true)
29610   if not file then return end
29611   local f = open(file,"rb")
29612   if not f then return end
29613   if offset and offset > 0 then
29614     f:seek("set", offset)
29615   end
29616   local data = f:read(length or 'a')
29617   f:close()
29618   return escapehex(data)
29619 end
29620 ltxutils.filedump = filedump
29621 deprecated(l3kernel, "filedump", function(name, offset, length)
29622   local dump = filedump(name, tonumber(offset), tonumber(length))
29623   if dump then
29624     write(dump)
29625   end
29626 end)

```

(End definition for ltx.utils.filedump and l3kernel.filedump. These functions are documented on page 260.)

`md5.HEX` Hash a string and return the hash in uppercase hexadecimal format. In some engines, this is build-in. For traditional LuaTeX, the conversion to hexadecimal has to be done by us.

```

29627 local md5_HEX = md5.HEX
29628 if not md5_HEX then
29629   local md5_sum = md5.sum
29630   function md5_HEX(data)
29631     return escapehex(md5_sum(data))
29632   end
29633   md5.HEX = md5_HEX
29634 end

```

(End definition for md5.HEX. This function is documented on page ??.)

`ltx.utils.filemd5sum` Read an entire file and hash it: the hash function itself is a built-in. As Lua is byte-
`l3kernel.filemdfivesum` based there is no work needed here in terms of UTF-8 (see `pdftexcmds` and how it handles strings that have passed through LuaTeX). The file is read in binary mode so that no line ending normalisation occurs.

```

29635 local function filemd5sum(name)
29636   local file = kpse_find(name, "tex", true) if not file then return end
29637   local f = open(file, "rb") if not f then return end
29638
29639   local data = f:read("*a")
29640   f:close()
29641   return md5_HEX(data)
29642 end
29643 ltxutils.filemd5sum = filemd5sum
29644 deprecated(l3kernel, "filemdfivesum", function(name)
29645   local hash = filemd5sum(name)

```



```

29646     if hash then
29647         write(hash)
29648     end
29649 end)

```

(End definition for `ltx.utils.filemd5sum` and `l3kernel.filemdfivesum`. These functions are documented on page 260.)

`ltx.utils.filemoddate` There are two cases: If the C standard library is C99 compliant, we can use `%z` to get the timezone in almost the right format. We only have to add primes and replace a zero or missing offset with Z.

Of course this would be boring, so Windows does things differently. There we have to manually calculate the offset. See procedure `makepdfmtime` in `utils.c` of pdfTeX.

```

29650 local filemoddate
29651 if os_date'%z':match'^[+-]%d%d%d$' then
29652     local pattern = lpeg.Cs(16 *
29653         (lpeg.Cg(lpeg.S'+-' * '0000' * lpeg.Cc'Z')
29654         + 3 * lpeg.Cc'"'" * 2 * lpeg.Cc'"'"
29655         + lpeg.Cc'Z')
29656     * -1)
29657 function filemoddate(name)
29658     local file = kpse_find(name, "tex", true)
29659     if not file then return end
29660     local date = lfs_attr(file, "modification")
29661     if not date then return end
29662     return pattern:match(os_date("D:%Y%m%d%H%M%S%z", date))
29663 end
29664 else
29665     local function filemoddate(name)
29666         local file = kpse_find(name, "tex", true)
29667         if not file then return end
29668         local date = lfs_attr(file, "modification")
29669         if not date then return end
29670         local d = os_date("*t", date)
29671         local u = os_date("!*t", date)
29672         local off = 60 * (d.hour - u.hour) + d.min - u.min
29673         if d.year ~= u.year then
29674             if d.year > u.year then
29675                 off = off + 1440
29676             else
29677                 off = off - 1440
29678             end
29679         elseif d.yday ~= u.yday then
29680             if d.yday > u.yday then
29681                 off = off + 1440
29682             else
29683                 off = off - 1440
29684             end
29685         end
29686         local timezone
29687         if off == 0 then
29688             timezone = "Z"
29689         else
29690             if off < 0 then

```

```

29691         timezone = "-"
29692         off = -off
29693     else
29694         timezone = "+"
29695     end
29696     timezone = format("%s%02d'%02d'", timezone, hours // 60, hours % 60)
29697 end
29698 return format("D:%04d%02d%02d%02d%02d%s",
29699             d.year, d.month, d.day, d.hour, d.min, d.sec, timezone)
29700 end
29701 end
29702 ltxutils.filemoddate = filemoddate
29703 deprecated(l3kernel, "filemoddate", function(name)
29704     local hash = filemoddate(name)
29705     if hash then
29706         write(hash)
29707     end
29708 end)

```

(End definition for `ltx.utils.filemoddate` and `l3kernel.filemoddate`. These functions are documented on page 260.)

`ltx.utils.filesize`
`l3kernel.filesize`

A simple disk lookup.

```

29709 local function filesize(name)
29710     local file = kpse_find(name, "tex", true)
29711     if file then
29712         local size = lfs_attr(file, "size")
29713         if size then
29714             return size
29715         end
29716     end
29717 end
29718 ltxutils.filesize = filesize
29719 deprecated(l3kernel, "filesize", function(name)
29720     local size = filesize(name)
29721     if size then
29722         write(size)
29723     end
29724 end)

```

(End definition for `ltx.utils.filesize` and `l3kernel.filesize`. These functions are documented on page 260.)

`l3kernel.strcmp`

String comparison which gives the same results as pdfTeX's `\pdfstrcmp`, although the ordering should likely not be relied upon!

```

29725 deprecated(l3kernel, "strcmp", function (A, B)
29726     if A == B then
29727         write("0")
29728     elseif A < B then
29729         write("-1")
29730     else
29731         write("1")
29732     end
29733 end)

```

(End definition for `l3kernel.strcmp`. This function is documented on page 261.)

`l3kernel.shellescape` Replicating the pdfTeX log interaction for shell escape.

```

29734 local os_exec      = os.execute
29735 deprecated(l3kernel, "shellescape", function(cmd)
29736   local status,msg = os_exec(cmd)
29737   if status == nil then
29738     write_nl("log","runsystem(" .. cmd .. ")...(" .. msg .. ")\n")
29739   elseif status == 0 then
29740     write_nl("log","runsystem(" .. cmd .. ")...executed\n")
29741   else
29742     write_nl("log","runsystem(" .. cmd .. ")...failed " .. (msg or "") .. "\n")
29743   end
29744 end)

```

(End definition for `l3kernel.shellescape`. This function is documented on page 261.)

`luaedef` An internal function for defining control sequences from Lua which behave like primitives. This acts as a wrapper around `token.set_lua` which accepts a function instead of an index into the functions table.

```

29745 local luacmd do
29746   local token_create = token.create
29747   local set_lua = token.set_lua
29748   local undefined_cs = token.command_id'undefined_cs'
29749
29750   if not context and not luatexbase then require'ltluatex' end
29751   if luatexbase then
29752     local new_luafunction = luatexbase.new_luafunction
29753     local functions = lua.get_functions_table()
29754     function luacmd(name, func, ...)
29755       local id
29756       local tok = token_create(name)
29757       if tok.command == undefined_cs then
29758         id = new_luafunction(name)
29759         set_lua(name, id, ...)
29760       else
29761         id = tok.index or tok.mode
29762       end
29763       functions[id] = func
29764     end
29765   elseif context then
29766     local register = context.functions.register
29767     local functions = context.functions.known
29768     function luacmd(name, func, ...)
29769       local tok = token.create(name)
29770       if tok.command == undefined_cs then
29771         token.set_lua(name, register(func), ...)
29772       else
29773         functions[tok.index or tok.mode] = func
29774       end
29775     end
29776   end
29777 end

```

(End definition for `lua def.`)

```
29778 </lua>
29779 </package>
```

46 l3unicode implementation

```
29780 (*package)
29781 @@=char
```

Case changing both for strings and “text” requires data from the Unicode Consortium. Some of this is build in to the format (as `\lccode` and `\uccode` values) but this covers only the simple one-to-one situations and does not fully handle for example case folding.

As only the data needs to remain at the end of this process, everything is set up inside a group. The only thing that is outside is creating a stream: they are global anyway and it is best to force a stream for all engines. For performance reasons, some of the code here is very low-level: the material is read during loading `expl3` in package mode.

```
29782 \ior_new:N \g__char_data_ior
29783 \bool_lazy_or:nnTF { \sys_if_engine_luatex_p: } { \sys_if_engine_xetex_p: }
29784 {
29785   \group_begin:
```

Access the primitive but suppress further expansion: active chars are otherwise an issue.

```
29786   \cs_set:Npn \__char_generate_char:n #1
29787   { \tex_detokenize:D \tex_expandafter:D { \tex_Uchar:D " #1 } }
```

A fast local implementation for generating characters; the chars may be active, so we prevent further expansion.

```
29788   \cs_set:Npx \__char_generate:n #1
29789   {
29790     \exp_not:N \tex_unexpanded:D \exp_not:N \exp_after:wN
29791     {
29792       \exp_not:N \tex_Ucharcat:D
29793       #1 ~
29794       \tex_catcode:D #1 ~
29795     }
29796   }
```

Parse the main Unicode data file for two things. First, we want the titlecase exceptions: the one-to-one lower- and uppercase mappings it contains are all be covered by the `TEX` data. Second, we need normalization data: at present, just the canonical NFD mappings. Those all yield either one or two codepoints, so the split is relatively easy.

```
29797   \ior_open:Nn \g__char_data_ior { UnicodeData.txt }
29798   \cs_set_protected:Npn \__char_data_auxi:w
29799   #1 ; #2 ; #3 ; #4 ; #5 ; #6 ; #7 ; #8 ; #9 ;
29800   {
29801     \tl_if_blank:nF {#6}
29802     {
29803       \tl_if_head_eq_charcode:nNF {#6} < % >
29804       { \__char_data_auxiii:w #1 ; #6 ~ \q_stop }
29805     }
29806     \__char_data_auxiii:w #1 ;
```

```

29807     }
29808 \cs_set_protected:Npn \__char_data_auxii:w #1 ; #2 ~ #3 \q_stop
29809 {
29810     \tl_const:cx
29811     { c__char_nfd_ \__char_generate_char:n {#1} _tl }
29812     {
29813         \__char_generate:n { "#2 }
29814         \tl_if_blank:nF {#3}
29815         { \__char_generate:n { "#3 } }
29816     }
29817 }
29818 \cs_set_protected:Npn \__char_data_auxiii:w
29819 #1 ; #2 ; #3 ; #4 ; #5 ; #6 ; #7 ~ \q_stop
29820 {
29821     \cs_set_nopar:Npn \l__char_tmpa_tl {#7}
29822     \reverse_if:N \if_meaning:w \l__char_tmpa_tl \c_empty_tl
29823     \cs_set_nopar:Npn \l__char_tmpb_tl {#5}
29824     \reverse_if:N \if_meaning:w \l__char_tmpa_tl \l__char_tmpb_tl
29825     \tl_const:cx
29826     { c__char_titlecase_ \__char_generate_char:n {#1} _tl }
29827     { \__char_generate:n { "#7 } }
29828     \fi:
29829     \fi:
29830 }
29831 \group_begin:
29832 \char_set_catcode_space:n { ‘ \ }%
29833 \ior_map_variable:NNn \g__char_data_ior \l__char_tmpa_tl
29834 {%
29835     \if_meaning:w \l__char_tmpa_tl \c_space_tl
29836     \exp_after:wN \ior_map_break:
29837     \fi:
29838     \exp_after:wN \__char_data_auxi:w \l__char_tmpa_tl \q_stop
29839 }%
29840 \group_end:
29841 \ior_close:N \g__char_data_ior

```

The other data files all use C-style comments so we have to worry about # tokens (and reading as strings). The set up for case folding is in two parts. For the basic (core) mappings, folding is the same as lower casing in most positions so only store the differences. For the more complex foldings, always store the result, splitting up the two or three code points in the input as required.

```

29842 \ior_open:Nn \g__char_data_ior { CaseFolding.txt }
29843 \cs_set_protected:Npn \__char_data_auxi:w #1 ;~ #2 ;~ #3 ; #4 \q_stop
29844 {
29845     \if:w \tl_head:n { #2 ? } C
29846     \reverse_if:N \if_int_compare:w
29847     \char_value_lccode:n {"#1} = "#3 ~
29848     \tl_const:cx
29849     { c__char_foldcase_ \__char_generate_char:n {#1} _tl }
29850     { \__char_generate:n { "#3 } }
29851     \fi:
29852     \else:
29853     \if:w \tl_head:n { #2 ? } F
29854     \__char_data_auxii:w #1 ~ #3 ~ \q_stop

```

```

29855         \fi:
29856     \fi:
29857 }
29858 \cs_set_protected:Npn \__char_data_auxii:w #1 ~ #2 ~ #3 ~ #4 \q_stop
29859 {
29860     \tl_const:cx { c__char_foldcase_ \__char_generate_char:n {#1} _tl }
29861     {
29862         \__char_generate:n { "#2 }
29863         \__char_generate:n { "#3 }
29864         \tl_if_blank:nF {#4}
29865         { \__char_generate:n { \int_value:w "#4 } }
29866     }
29867 }
29868 \ior_str_map_inline:Nn \g__char_data_ior
29869 {
29870     \reverse_if:N \if:w \c_hash_str \tl_head:w #1 \c_hash_str \q_stop
29871     \__char_data_auxi:w #1 \q_stop
29872     \fi:
29873 }
29874 \ior_close:N \g__char_data_ior

```

For upper- and lowercasing special situations, there is a bit more to do as we also have title casing to consider, plus we need to stop part-way through the file.

```

29875 \ior_open:Nn \g__char_data_ior { SpecialCasing.txt }
29876 \cs_set_protected:Npn \__char_data_auxi:w
29877 #1 ;~ #2 ;~ #3 ;~ #4 ; #5 \q_stop
29878 {
29879     \use:n { \__char_data_auxii:w #1 ~ lower ~ #2 ~ } ~ \q_stop
29880     \use:n { \__char_data_auxii:w #1 ~ upper ~ #4 ~ } ~ \q_stop
29881     \str_if_eq:nnF {#3} {#4}
29882     { \use:n { \__char_data_auxii:w #1 ~ title ~ #3 ~ } ~ \q_stop }
29883 }
29884 \cs_set_protected:Npn \__char_data_auxii:w
29885 #1 ~ #2 ~ #3 ~ #4 ~ #5 \q_stop
29886 {
29887     \tl_if_empty:nF {#4}
29888     {
29889         \tl_const:cx { c__char_ #2 case_ \__char_generate_char:n {#1} _tl }
29890         {
29891             \__char_generate:n { "#3 }
29892             \__char_generate:n { "#4 }
29893             \tl_if_blank:nF {#5}
29894             { \__char_generate:n { "#5 } }
29895         }
29896     }
29897 }
29898 \ior_str_map_inline:Nn \g__char_data_ior
29899 {
29900     \str_if_eq:eeTF
29901     { \tl_head:w #1 \c_hash_str \q_stop }
29902     { \c_hash_str }
29903     {
29904         \str_if_eq:eeT
29905         {#1}
29906         { \c_hash_str \c_space_tl Conditional~Mappings }

```

```

29907         { \ior_map_break: }
29908     }
29909     { \_char_data_auxi:w #1 \q_stop }
29910 }
29911 \ior_close:N \g__char_data_ior
29912 \group_end:
29913 }

```

For the 8-bit engines, the above is skipped but there is still some set up required. As case changing can only be applied to bytes, and they have to be in the ASCII range, we define a series of data stores to represent them, and the data are used such that only these are ever case-changed. We do open and close one file to force allocation of a read: this keeps all engines in line.

```

29914 {
29915     \group_begin:
29916     \cs_set_protected:Npn \_char_tmp:NN #1#2
29917     {
29918         \quark_if_recursion_tail_stop:N #2
29919         \tl_const:cn { c__char_uppercase_ #2 _tl } {#1}
29920         \tl_const:cn { c__char_lowercase_ #1 _tl } {#2}
29921         \tl_const:cn { c__char_foldcase_ #1 _tl } {#2}
29922         \_char_tmp:NN
29923     }
29924     \_char_tmp:NN
29925     AaBbCcDdEeFfGgHhIiJjKkLlMmNnOoPpQqRrSsTtUuVvWwXxYyZz
29926     ? \q_recursion_tail \q_recursion_stop
29927     \ior_open:Nn \g__char_data_ior { UnicodeData.txt }
29928     \ior_close:N \g__char_data_ior
29929     \group_end:
29930 }
29931 \</package>

```

47 l3text implementation

```

29932 \*package>
29933 \@@=text>

```

47.1 Internal auxiliaries

\s__text_stop Internal scan marks.

```

29934 \scan_new:N \s__text_stop

```

(End definition for \s__text_stop.)

\q__text_nil Internal quarks.

```

29935 \quark_new:N \q__text_nil

```

(End definition for \q__text_nil.)

__text_quark_if_nil_p:n Branching quark conditional.

```

\__text_quark_if_nil:nTF 29936 \__kernel_quark_new_conditional:Nn \__text_quark_if_nil:n { TF }

```

(End definition for __text_quark_if_nil:nTF.)

```

\q__text_recursion_tail Internal recursion quarks.
\q__text_recursion_stop 29937 \quark_new:N \q__text_recursion_tail
29938 \quark_new:N \q__text_recursion_stop

(End definition for \q__text_recursion_tail and \q__text_recursion_stop.)

__text_use_i_delimit_by_q_recursion_stop:nw Functions to gobble up to a quark.
29939 \cs_new:Npn \__text_use_i_delimit_by_q_recursion_stop:nw
29940 #1 #2 \q__text_recursion_stop {#1}

(End definition for \__text_use_i_delimit_by_q_recursion_stop:nw.)

\__text_if_recursion_tail_stop_do:Nn Functions to query recursion quarks.
29941 \__kernel_quark_new_test:N \__text_if_recursion_tail_stop_do:Nn

(End definition for \__text_if_recursion_tail_stop_do:Nn.)

```

47.2 Utilities

```

\__text_token_to_explicit:N The idea here is to take a token and ensure that if it's an implicit char, we output the
\__text_token_to_explicit_char:N explicit version. Otherwise, the token needs to be unchanged. First, we have to split
\__text_token_to_explicit_cs:N between control sequences and everything else.
\__text_token_to_explicit_cs_aux:N 29942 \group_begin:
\__text_token_to_explicit:n 29943 \char_set_catcode_active:n { 0 }
\__text_token_to_explicit_auxi:w 29944 \cs_new:Npn \__text_token_to_explicit:N #1
\__text_token_to_explicit_auxii:w 29945 {
\__text_token_to_explicit_auxiii:w 29946 \if_catcode:w \exp_not:N #1
29947 \if_catcode:w \scan_stop: \exp_not:N #1
29948 \scan_stop:
29949 \else:
29950 \exp_not:N ^^@
29951 \fi:
29952 \exp_after:wN \__text_token_to_explicit_cs:N
29953 \else:
29954 \exp_after:wN \__text_token_to_explicit_char:N
29955 \fi:
29956 #1
29957 }
29958 \group_end:

```

For control sequences, we can check for macros versus other cases using `\if_meaning:w`, then explicitly check for `\chardef` and `\mathchardef`.

```

29959 \cs_new:Npn \__text_token_to_explicit_cs:N #1
29960 {
29961 \exp_after:wN \if_meaning:w \exp_not:N #1 #1
29962 \exp_after:wN \use:nn \exp_after:wN
29963 \__text_token_to_explicit_cs_aux:N
29964 \else:
29965 \exp_after:wN \exp_not:n
29966 \fi:
29967 {#1}
29968 }
29969 \cs_new:Npn \__text_token_to_explicit_cs_aux:N #1
29970 {

```



```

29971 \bool_lazy_or:nnTF
29972 { \token_if_chardef_p:N #1 }
29973 { \token_if_mathchardef_p:N #1 }
29974 {
29975   \char_generate:nn {#1}
29976   { \char_value_catcode:n {#1} }
29977 }
29978 {#1}
29979 }

```

For character tokens, we need to filter out the implicit characters from those that are explicit. That's done here, then if necessary we work out the category code and generate the char. To avoid issues with alignment tabs, that one is done by elimination rather than looking up the code explicitly. The trick with finding the charcode is that the TeX messages are either the *⟨something⟩* character *⟨char⟩* or the *⟨type⟩* *⟨char⟩*.

```

29980 \cs_new:Npn \__text_token_to_explicit_char:N #1
29981 {
29982   \if:w
29983     \if_catcode:w ^ \exp_args:No \str_tail:n { \token_to_str:N #1 } ^
29984     \token_to_str:N #1 #1
29985     \else:
29986       AB
29987     \fi:
29988     \exp_after:wN \exp_not:n
29989   \else:
29990     \exp_after:wN \__text_token_to_explicit:n
29991   \fi:
29992   {#1}
29993 }
29994 \cs_new:Npn \__text_token_to_explicit:n #1
29995 {
29996   \exp_after:wN \__text_token_to_explicit_auxi:w
29997   \int_value:w
29998   \if_catcode:w \c_group_begin_token #1 1 \else:
29999   \if_catcode:w \c_group_end_token #1 2 \else:
30000   \if_catcode:w \c_math_toggle_token #1 3 \else:
30001   \if_catcode:w ## #1 6 \else:
30002   \if_catcode:w ^ #1 7 \else:
30003   \if_catcode:w \c_math_subscript_token #1 8 \else:
30004   \if_catcode:w \c_space_token #1 10 \else:
30005   \if_catcode:w A #1 11 \else:
30006   \if_catcode:w + #1 12 \else:
30007   4 \fi: \fi: \fi: \fi: \fi: \fi: \fi: \fi:
30008   \exp_after:wN ;
30009   \token_to_meaning:N #1 \s__text_stop
30010 }
30011 \cs_new:Npn \__text_token_to_explicit_auxi:w #1 ; #2 \s__text_stop
30012 {
30013   \char_generate:nn
30014   {
30015     \if_int_compare:w #1 < 9 \exp_stop_f:
30016     \exp_after:wN \__text_token_to_explicit_auxii:w
30017   \else:
30018     \exp_after:wN \__text_token_to_explicit_auxiii:w

```

```

30019         \fi:
30020         #2
30021     }
30022     {#1}
30023 }
30024 \exp_last_unbraced:NNNNo \cs_new:Npn \__text_token_to_explicit_auxii:w
30025   #1 { \tl_to_str:n { character ~ } } { ' }
30026 \cs_new:Npn \__text_token_to_explicit_auxiii:w #1 ~ #2 ~ { ' }

```

(End definition for `__text_token_to_explicit:N` and others.)

`__text_char_catcode:N` An idea from `l3char`: we need to get the category code of a specific token, not the general case.

```

30027 \cs_new:Npn \__text_char_catcode:N #1
30028 {
30029   \if_catcode:w \exp_not:N #1 \c_math_toggle_token
30030     3
30031   \else:
30032     \if_catcode:w \exp_not:N #1 \c_alignment_token
30033       4
30034     \else:
30035       \if_catcode:w \exp_not:N #1 \c_math_superscript_token
30036         7
30037     \else:
30038       \if_catcode:w \exp_not:N #1 \c_math_subscript_token
30039         8
30040     \else:
30041       \if_catcode:w \exp_not:N #1 \c_space_token
30042         10
30043     \else:
30044       \if_catcode:w \exp_not:N #1 \c_catcode_letter_token
30045         11
30046     \else:
30047       \if_catcode:w \exp_not:N #1 \c_catcode_other_token
30048         12
30049     \else:
30050       13
30051     \fi:
30052   \fi:
30053   \fi:
30054   \fi:
30055   \fi:
30056   \fi:
30057   \fi:
30058 }

```

(End definition for `__text_char_catcode:N`.)

`__text_if_expandable:NTF` Test for tokens that make sense to expand here: that is more restrictive than the engine view.

```

30059 \prg_new_conditional:Npnn \__text_if_expandable:N #1 { T , F , TF }
30060 {
30061   \token_if_expandable:NTF #1
30062   {

```

```

30063     \bool_lazy_any:nTF
30064     {
30065         { \token_if_protected_macro_p:N      #1 }
30066         { \token_if_protected_long_macro_p:N #1 }
30067         { \token_if_eq_meaning_p:NN \q__text_recursion_tail #1 }
30068     }
30069     { \prg_return_false: }
30070     { \prg_return_true: }
30071 }
30072 { \prg_return_false: }
30073 }

```

(End definition for `__text_if_expandable:NTF`.)

47.3 Configuration variables

`\l_text_accents_tl` `\l_text_letterlike_tl` Special cases for accents and letter-like symbols, which in some cases will need to be converted further.

```

30074 \tl_new:N \l_text_accents_tl
30075 \tl_set:Nn \l_text_accents_tl
30076 { \‘ \’ \^ \~ \= \u \. \" \r \H \v \d \c \k \b \t }
30077 \tl_new:N \l_text_letterlike_tl
30078 \tl_set:Nn \l_text_letterlike_tl
30079 {
30080     \AA \aa
30081     \AE \ae
30082     \DH \dh
30083     \DJ \dj
30084     \IJ \ij
30085     \L \l
30086     \NG \ng
30087     \O \o
30088     \OE \oe
30089     \SS \ss
30090     \TH \th
30091 }

```

(End definition for `\l_text_accents_tl` and `\l_text_letterlike_tl`. These variables are documented on page 266.)

`\l_text_case_exclude_arg_tl` Non-text arguments.

```

30092 \tl_new:N \l_text_case_exclude_arg_tl
30093 \tl_set:Nn \l_text_case_exclude_arg_tl { \begin \cite \end \label \ref }

```

(End definition for `\l_text_case_exclude_arg_tl`. This variable is documented on page 266.)

`\l_text_math_arg_tl` Math mode as arguments.

```

30094 \tl_new:N \l_text_math_arg_tl
30095 \tl_set:Nn \l_text_math_arg_tl { \ensuremath }

```

(End definition for `\l_text_math_arg_tl`. This variable is documented on page 266.)

`\l_text_math_delims_tl` Paired math mode delimiters.

```

30096 \tl_new:N \l_text_math_delims_tl
30097 \tl_set:Nn \l_text_math_delims_tl { $ $ \ ( \ ) }

```

(End definition for `\l_text_math_delims_tl`. This variable is documented on page 266.)

`\l_text_expand_exclude_tl` Commands which need not to expand.

```
30098 \tl_new:N \l_text_expand_exclude_tl
30099 \tl_set:Nn \l_text_expand_exclude_tl
30100 { \begin \cite \end \label \ref }
```

(End definition for `\l_text_expand_exclude_tl`. This variable is documented on page 266.)

`\l__text_math_mode_tl` Used to control math mode output: internal as there is a dedicated setter.

```
30101 \tl_new:N \l__text_math_mode_tl
```

(End definition for `\l__text_math_mode_tl`.)

47.4 Expansion to formatted text

Markers for implicit char handling.

```
\c__text_chardef_space_token
\c__text_mathchardef_space_token
\c__text_chardef_group_begin_token
\c__text_mathchardef_group_begin_token
\c__text_chardef_group_end_token
\c__text_mathchardef_group_end_token
30102 \tex_chardef:D \c__text_chardef_space_token = '\ %
30103 \tex_mathchardef:D \c__text_mathchardef_space_token = '\ %
30104 \tex_chardef:D \c__text_chardef_group_begin_token = '\{ % \}
30105 \tex_mathchardef:D \c__text_mathchardef_group_begin_token = '\{ % \} \{
30106 \tex_chardef:D \c__text_chardef_group_end_token = '\} % \{
30107 \tex_mathchardef:D \c__text_mathchardef_group_end_token = '\} %
```

(End definition for `\c__text_chardef_space_token` and others.)

`\text_expand:n`

After precautions against & tokens, start a simple loop: that of course means that “text” cannot contain the two recursion quarks. The loop here must be f-type expandable; we have arbitrary user commands which might be protected *and* take arguments, and if the expansion code is used in a typesetting context, that will otherwise explode. (The same issue applies more clearly to case changing: see the example there.)

```
30108 \cs_new:Npn \text_expand:n #1
30109 {
30110   \__kernel_exp_not:w \exp_after:wN
30111   {
30112     \exp:w
30113     \__text_expand:n {#1}
30114   }
30115 }
30116 \cs_new:Npn \__text_expand:n #1
30117 {
30118   \group_align_safe_begin:
30119   \__text_expand_loop:w #1
30120   \q__text_recursion_tail \q__text_recursion_stop
30121   \__text_expand_result:n { }
30122 }
```

The approach to making the code f-type expandable is to use a marker result token and to shuffle the collected tokens

```
30123 \cs_new:Npn \__text_expand_store:n #1
30124 { \__text_expand_store:nw {#1} }
30125 \cs_generate_variant:Nn \__text_expand_store:n { o }
30126 \cs_new:Npn \__text_expand_store:nw #1#2 \__text_expand_result:n #3
30127 { #2 \__text_expand_result:n { #3 #1 } }
```

```

30128 \cs_new:Npn \__text_expand_end:w #1 \__text_expand_result:n #2
30129 {
30130   \group_align_safe_end:
30131   \exp_end:
30132   #2
30133 }

```

The main loop is a standard “tl action”; groups are handled recursively, while spaces are just passed through. Thus all of the action is in handling N-type tokens.

```

30134 \cs_new:Npn \__text_expand_loop:w #1 \q_text_recursion_stop
30135 {
30136   \tl_if_head_is_N_type:nTF {#1}
30137   { \__text_expand_N_type:N }
30138   {
30139     \tl_if_head_is_group:nTF {#1}
30140     { \__text_expand_group:n }
30141     { \__text_expand_space:w }
30142   }
30143   #1 \q_text_recursion_stop
30144 }
30145 \cs_new:Npn \__text_expand_group:n #1
30146 {
30147   \__text_expand_store:o
30148   {
30149     \exp_after:wN
30150     {
30151       \exp:w
30152       \__text_expand:n {#1}
30153     }
30154   }
30155   \__text_expand_loop:w
30156 }
30157 \exp_last_unbraced:NNo \cs_new:Npn \__text_expand_space:w \c_space_tl
30158 {
30159   \__text_expand_store:n { ~ }
30160   \__text_expand_loop:w
30161 }

```

Before we get into the real work, we have to watch out for problematic implicit characters: spaces and grouping tokens. Converting these to explicit characters later would lead to real issues as they are *not* N-type. A space is the easy case, so it’s dealt with first: just insert the explicit token and continue the loop.

```

30162 \cs_new:Npx \__text_expand_N_type:N #1
30163 {
30164   \exp_not:N \__text_if_recursion_tail_stop_do:Nn #1
30165   { \exp_not:N \__text_expand_end:w }
30166   \exp_not:N \bool_lazy_any:nTF
30167   {
30168     { \exp_not:N \token_if_eq_meaning_p:NN #1 \c_space_token }
30169     {
30170       \exp_not:N \token_if_eq_meaning_p:NN #1
30171       \c_text_chardef_space_token
30172     }
30173   }
30174   \exp_not:N \token_if_eq_meaning_p:NN #1

```

```

30175         \c__text_mathchardef_space_token
30176     }
30177 }
30178 { \exp_not:N \__text_expand_space:w \c_space_tl }
30179 { \exp_not:N \__text_expand_N_type_auxi:N #1 }
30180 }

```

Implicit `{/}` offer two issues. First, the token could be an implicit brace character: we need to avoid turning that into a brace group, so filter out the cases manually. Then we handle the case where an implicit group is present. That is done in an “open-ended” way: there’s the possibility the closing token is hidden somewhere.

```

30181 \cs_new:Npn \__text_expand_N_type_auxi:N #1
30182 {
30183     \bool_lazy_or:nnTF
30184     { \token_if_eq_meaning_p:NN #1 \c__text_chardef_group_begin_token }
30185     { \token_if_eq_meaning_p:NN #1 \c__text_mathchardef_group_begin_token }
30186     {
30187         \__text_expand_store:o \c_left_brace_str
30188         \__text_expand_loop:w
30189     }
30190     {
30191         \bool_lazy_or:nnTF
30192         { \token_if_eq_meaning_p:NN #1 \c__text_chardef_group_end_token }
30193         { \token_if_eq_meaning_p:NN #1 \c__text_mathchardef_group_end_token }
30194         {
30195             \__text_expand_store:o \c_right_brace_str
30196             \__text_expand_loop:w
30197         }
30198         { \__text_expand_N_type_auxii:N #1 }
30199     }
30200 }
30201 \cs_new:Npn \__text_expand_N_type_auxii:N #1
30202 {
30203     \token_if_eq_meaning:NNTF #1 \c_group_begin_token
30204     {
30205         { \if_false: } \fi:
30206         \__text_expand_loop:w
30207     }
30208     {
30209         \token_if_eq_meaning:NNTF #1 \c_group_end_token
30210         {
30211             \if_false: { \fi: }
30212             \__text_expand_loop:w
30213         }
30214         { \__text_expand_N_type_auxiii:N #1 }
30215     }
30216 }

```

The first step in dealing with N-type tokens is to look for math mode material: that needs to be left alone. The starting function has to be split into two as we need `\quark_if_recursion_tail_stop:N` first before we can trigger the search. We then look for matching pairs of delimiters, allowing for the case where math mode starts but does not end. Within math mode, we simply pass all the tokens through unchanged, just checking the N-type ones against the end marker.

```

30217 \cs_new:Npn \__text_expand_N_type_auxiii:N #1
30218 {
30219   \exp_after:wN \__text_expand_math_search:NNN
30220   \exp_after:wN #1 \l_text_math_delims_tl
30221   \q__text_recursion_tail \q__text_recursion_tail
30222   \q__text_recursion_stop
30223 }
30224 \cs_new:Npn \__text_expand_math_search:NNN #1#2#3
30225 {
30226   \__text_if_recursion_tail_stop_do:Nn #2
30227   { \__text_expand_explicit:N #1 }
30228   \token_if_eq_meaning:NNTF #1 #2
30229   {
30230     \__text_use_i_delimit_by_q_recursion_stop:nw
30231     {
30232       \__text_expand_store:n {#1}
30233       \__text_expand_math_loop:Nw #3
30234     }
30235   }
30236   { \__text_expand_math_search:NNN #1 }
30237 }
30238 \cs_new:Npn \__text_expand_math_loop:Nw #1#2 \q__text_recursion_stop
30239 {
30240   \tl_if_head_is_N_type:nTF {#2}
30241   { \__text_expand_math_N_type:NN }
30242   {
30243     \tl_if_head_is_group:nTF {#2}
30244     { \__text_expand_math_group:Nn }
30245     { \__text_expand_math_space:Nw }
30246   }
30247   #1#2 \q__text_recursion_stop
30248 }
30249 \cs_new:Npn \__text_expand_math_N_type:NN #1#2
30250 {
30251   \__text_if_recursion_tail_stop_do:Nn #2
30252   { \__text_expand_end:w }
30253   \__text_expand_store:n {#2}
30254   \token_if_eq_meaning:NNTF #2 #1
30255   { \__text_expand_loop:w }
30256   { \__text_expand_math_loop:Nw #1 }
30257 }
30258 \cs_new:Npn \__text_expand_math_group:Nn #1#2
30259 {
30260   \__text_expand_store:n { {#2} }
30261   \__text_expand_math_loop:Nw #1
30262 }
30263 \exp_after:wN \cs_new:Npn \exp_after:wN \__text_expand_math_space:Nw
30264 \exp_after:wN # \exp_after:wN 1 \c_space_tl
30265 {
30266   \__text_expand_store:n { ~ }
30267   \__text_expand_math_loop:Nw #1
30268 }

```

At this stage, either we have a control sequence or a simple character: split and handle.

```

30269 \cs_new:Npn \__text_expand_explicit:N #1
30270 {
30271   \token_if_cs:NTF #1
30272   { \__text_expand_exclude:N #1 }
30273   {
30274     \__text_expand_store:n {#1}
30275     \__text_expand_loop:w
30276   }
30277 }

```

Next we exclude math commands: this is mainly as there *might* be an `\ensuremath`. We also handle accents, which are basically the same issue but are kept separate for semantic reasons.

```

30278 \cs_new:Npn \__text_expand_exclude:N #1
30279 {
30280   \exp_args:Ne \__text_expand_exclude:nN
30281   {
30282     \exp_not:V \l_text_math_arg_tl
30283     \exp_not:V \l_text_accents_tl
30284     \exp_not:V \l_text_expand_exclude_tl
30285   }
30286   #1
30287 }
30288 \cs_new:Npn \__text_expand_exclude:nN #1#2
30289 {
30290   \__text_expand_exclude:NN #2 #1
30291   \q__text_recursion_tail \q__text_recursion_stop
30292 }
30293 \cs_new:Npn \__text_expand_exclude:NN #1#2
30294 {
30295   \__text_if_recursion_tail_stop_do:Nn #2
30296   { \__text_expand_letterlike:N #1 }
30297   \str_if_eq:nnTF {#1} {#2}
30298   {
30299     \__text_use_i_delimit_by_q_recursion_stop:nw
30300     { \__text_expand_exclude:Nn #1 }
30301   }
30302   { \__text_expand_exclude:NN #1 }
30303 }
30304 \cs_new:Npn \__text_expand_exclude:Nn #1#2
30305 {
30306   \__text_expand_store:n { #1 {#2} }
30307   \__text_expand_loop:w
30308 }

```

Another list of exceptions: these ones take no arguments so are easier to handle.

```

30309 \cs_new:Npn \__text_expand_letterlike:N #1
30310 {
30311   \exp_after:wN \__text_expand_letterlike:NN \exp_after:wN
30312   #1 \l_text_letterlike_tl
30313   \q__text_recursion_tail \q__text_recursion_stop
30314 }
30315 \cs_new:Npn \__text_expand_letterlike:NN #1#2
30316 {
30317   \__text_if_recursion_tail_stop_do:Nn #2

```



```

30318     { \_text_expand_cs:N #1 }
30319 \cs_if_eq:NNTF #2 #1
30320     {
30321         \_text_use_i_delimit_by_q_recursion_stop:nw
30322         {
30323             \_text_expand_store:n {#1}
30324             \_text_expand_loop:w
30325         }
30326     }
30327     { \_text_expand_letterlike:NN #1 }
30328 }

```

LaTeX 2_ε's `\protect` makes life interesting. Where possible, we simply remove it and replace with the “parent” command; of course, the `\protect` might be explicit, in which case we need to leave it alone if it's required. There is also the case of a straight `\@protected@testopt` to cover.

```

30329 \cs_new:Npx \_text_expand_cs:N #1
30330 {
30331     \exp_not:N \str_if_eq:nnTF {#1} { \exp_not:N \protect }
30332     { \exp_not:N \_text_expand_protect:N }
30333     {
30334         \bool_lazy_and:nnTF
30335         { \cs_if_exist_p:N \fmtname }
30336         { \str_if_eq_p:Vn \fmtname { LaTeX2e } }
30337         { \exp_not:N \_text_expand_testopt:N #1 }
30338         { \exp_not:N \_text_expand_replace:N #1 }
30339     }
30340 }
30341 \cs_new:Npn \_text_expand_protect:N #1
30342 {
30343     \exp_args:Ne \_text_expand_protect:nN
30344     { \cs_to_str:N #1 } #1
30345 }
30346 \cs_new:Npn \_text_expand_protect:nN #1#2
30347 { \_text_expand_protect:Nw #2 #1 \q__text_nil #1 ~ \q__text_nil \q__text_nil \s__text_stop
30348 \cs_new:Npn \_text_expand_protect:Nw #1 #2 ~ \q__text_nil #3 \q__text_nil #4 \s__text_stop
30349 {
30350     \_text_quark_if_nil:nTF {#4}
30351     {
30352         \cs_if_exist:cTF {#2}
30353         { \exp_args:Ne \_text_expand_store:n { \exp_not:c {#2} } }
30354         { \_text_expand_store:n { \protect #1 } }
30355     }
30356     { \_text_expand_store:n { \protect #1 } }
30357     \_text_expand_loop:w
30358 }
30359 \cs_new:Npn \_text_expand_testopt:N #1
30360 {
30361     \str_if_eq:nnTF {#1} { \@protected@testopt }
30362     { \_text_expand_testopt:NNn }
30363     { \_text_expand_encoding:N #1 }
30364 }
30365 \cs_new:Npn \_text_expand_testopt:NNn #1#2#3
30366 {

```

```

30367     \__text_expand_store:n {#1}
30368     \__text_expand_loop:w
30369 }

```

Deal with encoding-specific commands

```

30370 \cs_new:Npn \__text_expand_encoding:N #1
30371 {
30372     \bool_lazy_or:nnTF
30373     { \cs_if_eq_p:NN #1 \@current@cmd }
30374     { \cs_if_eq_p:NN #1 \@changed@cmd }
30375     { \exp_after:wN \__text_expand_loop:w \__text_expand_encoding_escape:NN }
30376     { \__text_expand_replace:N #1 }
30377 }
30378 \cs_new:Npn \__text_expand_encoding_escape:NN #1#2 { \exp_not:n {#1} }

```

See if there is a dedicated replacement, and if there is, insert it.

```

30379 \cs_new:Npn \__text_expand_replace:N #1
30380 {
30381     \bool_lazy_and:nnTF
30382     { \cs_if_exist_p:c { l__text_expand_ \token_to_str:N #1 _t1 } }
30383     {
30384         \bool_lazy_or_p:nn
30385         { \token_if_cs_p:N #1 }
30386         { \token_if_active_p:N #1 }
30387     }
30388     {
30389         \exp_args:Nv \__text_expand_replace:n
30390         { l__text_expand_ \token_to_str:N #1 _t1 }
30391     }
30392     { \__text_expand_cs_expand:N #1 }
30393 }
30394 \cs_new:Npn \__text_expand_replace:n #1 { \__text_expand_loop:w #1 }

```

Finally, expand any macros which can be: this then loops back around to deal with what they produce. The only issue is if the token is `\exp_not:n`, as that must apply to the following balanced text. There might be an `\exp_after:wN` there, so we check for it.

```

30395 \cs_new:Npn \__text_expand_cs_expand:N #1
30396 {
30397     \__text_if_expandable:NTF #1
30398     {
30399         \token_if_eq_meaning:NNTF #1 \exp_not:n
30400         { \__text_expand_noexpand:w }
30401         { \exp_after:wN \__text_expand_loop:w #1 }
30402     }
30403     {
30404         \__text_expand_store:n {#1}
30405         \__text_expand_loop:w
30406     }
30407 }
30408 \cs_new:Npn \__text_expand_noexpand:w #1#
30409 { \__text_expand_noexpand:nn {#1} }
30410 \cs_new:Npn \__text_expand_noexpand:nn #1#2
30411 {
30412     #1 \__text_expand_store:n #1 {#2}
30413     \__text_expand_loop:w

```

```
30414 }
```

(End definition for `\text_expand:n` and others. This function is documented on page 263.)

```
\text_declare_expand_equivalent:Nn
\text_declare_expand_equivalent:cn
```

Create equivalents to allow replacement.

```
30415 \cs_new_protected:Npn \text_declare_expand_equivalent:Nn #1#2
30416 {
30417   \tl_clear_new:c { l__text_expand_ \token_to_str:N #1 _tl }
30418   \tl_set:cn { l__text_expand_ \token_to_str:N #1 _tl } {#2}
30419 }
30420 \cs_generate_variant:Nn \text_declare_expand_equivalent:Nn { c }
```

(End definition for `\text_declare_expand_equivalent:Nn`. This function is documented on page 263.)

```
30421 \</package>
```

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```
30422 \*package>
```

```
30423 \@@=text>
```

48.1 Case changing

```
\l_text_titlecase_check_letter_bool
```

Needed to determine the route used in titlecasing.

```
30424 \bool_new:N \l_text_titlecase_check_letter_bool
30425 \bool_set_true:N \l_text_titlecase_check_letter_bool
```

(End definition for `\l_text_titlecase_check_letter_bool`. This variable is documented on page 266.)

```
\text_lowercase:n
\text_uppercase:n
\text_titlecase:n
\text_titlecase_first:n
\text_lowercase:nn
\text_uppercase:nn
\text_titlecase:nn
\text_titlecase_first:nn
```

The user level functions here are all wrappers around the internal functions for case changing.

```
30426 \cs_new:Npn \text_lowercase:n #1
30427 { \__text_change_case:nnn { lower } { } {#1} }
30428 \cs_new:Npn \text_uppercase:n #1
30429 { \__text_change_case:nnn { upper } { } {#1} }
30430 \cs_new:Npn \text_titlecase:n #1
30431 { \__text_change_case:nnn { title } { } {#1} }
30432 \cs_new:Npn \text_titlecase_first:n #1
30433 { \__text_change_case:nnn { titleonly } { } {#1} }
30434 \cs_new:Npn \text_lowercase:nn #1#2
30435 { \__text_change_case:nnn { lower } {#1} {#2} }
30436 \cs_new:Npn \text_uppercase:nn #1#2
30437 { \__text_change_case:nnn { upper } {#1} {#2} }
30438 \cs_new:Npn \text_titlecase:nn #1#2
30439 { \__text_change_case:nnn { title } {#1} {#2} }
30440 \cs_new:Npn \text_titlecase_first:nn #1#2
30441 { \__text_change_case:nnn { titleonly } {#1} {#2} }
```

(End definition for `\text_lowercase:n` and others. These functions are documented on page 265.)

As for the expansion code, the business end of case changing is the handling of N-type tokens. First, we expand the input fully (so the loops here don't need to worry about awkward look-aheads and the like). Then we split into the different paths.

The code here needs to be f-type expandable to deal with the situation where case changing is applied in running text. There, we might have case changing as a document command and the text containing other non-expandable document commands.

```
\__text_change_case:nnn
\__text_change_case_aux:nnn
\__text_change_case_store:n
\__text_change_case_store:o
\__text_change_case_store:V
\__text_change_case_store:v
\__text_change_case_store:e
\__text_change_case_store:nw
\__text_change_case_result:n
\__text_change_case_end:w
\__text_change_case_loop:nnw
\__text_change_case_break:w
\__text_change_case_group_lower:nnn
\__text_change_case_group_upper:nnn
\__text_change_case_group_title:nnn
\__text_change_case_group_titleonly:nnn
```

```

\cs_set_eq:NN \MakeLowercase \text_lowercase
...
\MakeLowercase{\enquote*{A} text}

```

If we use an e-type expansion and wrap each token in `\exp_not:n`, that would explode: the document command grabs `\exp_not:n` as an argument, and things go badly wrong. So we have to wrap the entire result in exactly one `\exp_not:n`, or rather in the kernel version.

```

30442 \cs_new:Npn \__text_change_case:nnn #1#2#3
30443 {
30444   \__kernel_exp_not:w \exp_after:wN
30445   {
30446     \exp:w
30447     \exp_args:Ne \__text_change_case_aux:nnn
30448     { \text_expand:n {#3} }
30449     {#1} {#2}
30450   }
30451 }
30452 \cs_new:Npn \__text_change_case_aux:nnn #1#2#3
30453 {
30454   \group_align_safe_begin:
30455   \cs_if_exist_use:c { __text_change_case_boundary_ #2 _ #3 :Nnnw }
30456   \__text_change_case_loop:nnw {#2} {#3} #1
30457   \q__text_recursion_tail \q__text_recursion_stop
30458   \__text_change_case_result:n { }
30459 }

```

As for expansion, collect up the tokens for future use.

```

30460 \cs_new:Npn \__text_change_case_store:n #1
30461 { \__text_change_case_store:nw {#1} }
30462 \cs_generate_variant:Nn \__text_change_case_store:n { o , e , V , v }
30463 \cs_new:Npn \__text_change_case_store:nw #1#2 \__text_change_case_result:n #3
30464 { #2 \__text_change_case_result:n { #3 #1 } }
30465 \cs_new:Npn \__text_change_case_end:w #1 \__text_change_case_result:n #2
30466 {
30467   \group_align_safe_end:
30468   \exp_end:
30469   #2
30470 }

```

The main loop is the standard `tl` action type.

```

30471 \cs_new:Npn \__text_change_case_loop:nnw #1#2#3 \q__text_recursion_stop
30472 {
30473   \tl_if_head_is_N_type:nTF {#3}
30474   { \__text_change_case_N_type:nnN }
30475   {
30476     \tl_if_head_is_group:nTF {#3}
30477     { \use:c { __text_change_case_group_ #1 :nnn } }
30478     { \__text_change_case_space:nnw }
30479   }
30480   {#1} {#2} #3 \q__text_recursion_stop
30481 }
30482 \cs_new:Npn \__text_change_case_break:w #1 \q__text_recursion_tail \q__text_recursion_stop
30483 {

```

```

30484     \__text_change_case_store:n {#1}
30485     \__text_change_case_end:w
30486 }

```

For a group, we *could* worry about whether this contains a character or not. However, that would make life very complex for little gain: exactly what a first character is is rather weakly-defined anyway. So if there is a group, we simply assume that a character has been seen, and for title case we switch to the “rest of the tokens” situation. To avoid having too much testing, we use a two-step process here to allow the titlecase functions to be separate.

```

30487 \cs_new:Npn \__text_change_case_group_lower:nnn #1#2#3
30488 {
30489     \__text_change_case_store:o
30490     {
30491         \exp_after:wN
30492         {
30493             \exp:w
30494             \__text_change_case_aux:nnn {#3} {#1} {#2}
30495         }
30496     }
30497     \__text_change_case_loop:nnw {#1} {#2}
30498 }
30499 \cs_new_eq:NN \__text_change_case_group_upper:nnn
30500 \__text_change_case_group_lower:nnn
30501 \cs_new:Npn \__text_change_case_group_title:nnn #1#2#3
30502 {
30503     \__text_change_case_store:o
30504     {
30505         \exp_after:wN
30506         {
30507             \exp:w
30508             \__text_change_case_aux:nnn {#3} {#1} {#2}
30509         }
30510     }
30511     \__text_change_case_loop:nnw { lower } {#2}
30512 }
30513 \cs_new:Npn \__text_change_case_group_titleonly:nnn #1#2#3
30514 {
30515     \__text_change_case_store:o
30516     {
30517         \exp_after:wN
30518         {
30519             \exp:w
30520             \__text_change_case_aux:nnn {#3} {#1} {#2}
30521         }
30522     }
30523     \__text_change_case_break:w
30524 }
30525 \use:x
30526 {
30527     \cs_new:Npn \exp_not:N \__text_change_case_space:nnw ##1##2 \c_space_tl
30528 }
30529 {
30530     \__text_change_case_store:n { ~ }

```

```

30531 \cs_if_exist_use:c { __text_change_case_boundary_ #1 _ #2 :Nnnw }
30532 \__text_change_case_loop:nnw {#1} {#2}
30533 }

```

The first step of handling N-type tokens is to filter out the end-of-loop. That has to be done separately from the first real step as otherwise we pick up the wrong delimiter. The loop here is the same as the `expand` one, just passing the additional data long. If no close-math token is found then the final clean-up is forced (i.e. there is no assumption of “well-behaved” input in terms of math mode).

```

30534 \cs_new:Npn \__text_change_case_N_type:nnN #1#2#3
30535 {
30536   \__text_if_recursion_tail_stop_do:Nn #3
30537   { \__text_change_case_end:w }
30538   \__text_change_case_N_type_aux:nnN {#1} {#2} #3
30539 }
30540 \cs_new:Npn \__text_change_case_N_type_aux:nnN #1#2#3
30541 {
30542   \exp_args:NV \__text_change_case_N_type:nnnN
30543   \l_text_math_delims_tl {#1} {#2} #3
30544 }
30545 \cs_new:Npn \__text_change_case_N_type:nnnN #1#2#3#4
30546 {
30547   \__text_change_case_math_search:nnNNN {#2} {#3} #4 #1
30548   \q__text_recursion_tail \q__text_recursion_tail
30549   \q__text_recursion_stop
30550 }
30551 \cs_new:Npn \__text_change_case_math_search:nnNNN #1#2#3#4#5
30552 {
30553   \__text_if_recursion_tail_stop_do:Nn #4
30554   { \__text_change_case_cs_check:nnN {#1} {#2} #3 }
30555   \token_if_eq_meaning:NNTF #3 #4
30556   {
30557     \__text_use_i_delimit_by_q_recursion_stop:nw
30558     {
30559       \__text_change_case_store:n {#3}
30560       \__text_change_case_math_loop:nnNw {#1} {#2} #5
30561     }
30562   }
30563   { \__text_change_case_math_search:nnNNN {#1} {#2} #3 }
30564 }
30565 \cs_new:Npn \__text_change_case_math_loop:nnNw #1#2#3#4 \q__text_recursion_stop
30566 {
30567   \tl_if_head_is_N_type:nTF {#4}
30568   { \__text_change_case_math_N_type:nnNN }
30569   {
30570     \tl_if_head_is_group:nTF {#4}
30571     { \__text_change_case_math_group:nnNn }
30572     { \__text_change_case_math_space:nnNw }
30573   }
30574   {#1} {#2} #3 #4 \q__text_recursion_stop
30575 }
30576 \cs_new:Npn \__text_change_case_math_N_type:nnNN #1#2#3#4
30577 {
30578   \__text_if_recursion_tail_stop_do:Nn #4

```

```

30579     { \_text_change_case_end:w }
30580 \_text_change_case_store:n {#4}
30581 \token_if_eq_meaning:NNTF #4 #3
30582   { \_text_change_case_loop:nnw {#1} {#2} }
30583   { \_text_change_case_math_loop:nnNw {#1} {#2} #3 }
30584 }
30585 \cs_new:Npn \_text_change_case_math_group:nnNn #1#2#3#4
30586 {
30587   \_text_change_case_store:n { {#4} }
30588   \_text_change_case_math_loop:nnNw {#1} {#2} #3
30589 }
30590 \use:x
30591 {
30592   \cs_new:Npn \exp_not:N \_text_change_case_math_space:nnNw ##1##2##3
30593     \c_space_tl
30594 }
30595 {
30596   \_text_change_case_store:n { ~ }
30597   \_text_change_case_math_loop:nnNw {#1} {#2} #3
30598 }

```

Once potential math-mode cases are filtered out the next stage is to test if the token grabbed is a control sequence: the two routes the code may take are then very different.

```

30599 \cs_new:Npn \_text_change_case_cs_check:nnN #1#2#3
30600 {
30601   \token_if_cs:NNTF #3
30602   { \_text_change_case_exclude:nnN }
30603   { \use:c { \_text_change_case_char_ #1 :nnN } }
30604     {#1} {#2} #3
30605 }

```

To deal with a control sequence there is first a need to test if it is on the list which indicate that case changing should be skipped. That's done using a loop as for the other special cases. If a hit is found then the argument is grabbed and passed through as-is.

```

30606 \cs_new:Npn \_text_change_case_exclude:nnN #1#2#3
30607 {
30608   \exp_args:Ne \_text_change_case_exclude:nnnN
30609   {
30610     \exp_not:V \l_text_math_arg_tl
30611     \exp_not:V \l_text_case_exclude_arg_tl
30612   }
30613   {#1} {#2} #3
30614 }
30615 \cs_new:Npn \_text_change_case_exclude:nnnN #1#2#3#4
30616 {
30617   \_text_change_case_exclude:nnNN {#2} {#3} #4 #1
30618   \q__text_recursion_tail \q__text_recursion_stop
30619 }
30620 \cs_new:Npn \_text_change_case_exclude:nnNN #1#2#3#4
30621 {
30622   \_text_if_recursion_tail_stop_do:Nn #4
30623   { \use:c { \_text_change_case_letterlike_ #1 :nnN } {#1} {#2} #3 }
30624   \str_if_eq:nnTF {#3} {#4}
30625   {
30626     \_text_use_i_delimit_by_q_recursion_stop:nw

```

```

30627         { \_text_change_case_exclude:nnNn {#1} {#2} #3 }
30628     }
30629     { \_text_change_case_exclude:nnNN {#1} {#2} #3 }
30630 }
30631 \cs_new:Npn \_text_change_case_exclude:nnNn #1#2#3#4
30632 {
30633     \_text_change_case_store:n { #3 {#4} }
30634     \_text_change_case_loop:nnw {#1} {#2}
30635 }

```

Letter-like commands may still be present: they are set up using a simple lookup approach, so can easily be handled with no loop. If there is no hit, we are at the end of the process: we loop around. Letter-like chars are all available only in upper- and lowercase, so titlecasing maps to the uppercase version.

```

30636 \cs_new:Npn \_text_change_case_letterlike_lower:nnN #1#2#3
30637 { \_text_change_case_letterlike:nnnnN {#1} {#1} {#1} {#2} #3 }
30638 \cs_new_eq:NN \_text_change_case_letterlike_upper:nnN
30639     \_text_change_case_letterlike_lower:nnN
30640 \cs_new:Npn \_text_change_case_letterlike_title:nnN #1#2#3
30641 { \_text_change_case_letterlike:nnnnN { upper } { lower } {#1} {#2} #3 }
30642 \cs_new:Npn \_text_change_case_letterlike_titleonly:nnN #1#2#3
30643 { \_text_change_case_letterlike:nnnnN { upper } { end } {#1} {#2} #3 }
30644 \cs_new:Npn \_text_change_case_letterlike:nnnnN #1#2#3#4#5
30645 {
30646     \cs_if_exist:cTF { c__text_ #1 case_ \token_to_str:N #5 _tl }
30647     {
30648         \_text_change_case_store:v
30649         { c__text_ #1 case_ \token_to_str:N #5 _tl }
30650         \use:c { __text_change_case_char_next_ #2 :nn } {#2} {#4}
30651     }
30652     {
30653         \_text_change_case_store:n {#5}
30654         \cs_if_exist:cTF
30655         {
30656             c__text_
30657             \str_if_eq:nnTF {#1} { lower } { upper } { lower }
30658             case_ \token_to_str:N #5 _tl
30659         }
30660         { \use:c { __text_change_case_char_next_ #2 :nn } {#2} {#4} }
30661         { \_text_change_case_loop:nnw {#3} {#4} }
30662     }
30663 }

```

For upper- and lowercase changes, once we get to this stage there are only a couple of questions remaining: is there a language-specific mapping and is there the special case of a terminal sigma. If not, then we pass to a simple character mapping.

```

30664 \cs_new:Npn \_text_change_case_char_lower:nnN #1#2#3
30665 {
30666     \cs_if_exist_use:cF { __text_change_case_lower_ #2 :nnnN }
30667     { \_text_change_case_lower_sigma:nnnN }
30668     {#1} {#1} {#2} #3
30669 }
30670 \cs_new:Npn \_text_change_case_char_upper:nnN #1#2#3
30671 {

```



```

30672 \cs_if_exist_use:cF { __text_change_case_upper_ #2 :nnnN }
30673 { \__text_change_case_char:nnnN }
30674 {#1} {#1} {#2} #3
30675 }

```

If the current character is an uppercase sigma, the a check is made on the next item in the input. If it is N-type and not a control sequence then there is a look-ahead phase: the logic here is simply based on letters. The one exception is Dutch: see below.

```

30676 \bool_lazy_or:nnTF
30677 { \sys_if_engine luatex_p: }
30678 { \sys_if_engine xetex_p: }
30679 {
30680 \cs_new:Npn \__text_change_case_lower_sigma:nnnN #1#2#3#4
30681 {
30682 \int_compare:nNnTF { '#4 } = { "03A3 }
30683 { \__text_change_case_lower_sigma:nnNw {#2} {#3} #4 }
30684 { \__text_change_case_char:nnnN {#1} {#2} {#3} #4 }
30685 }
30686 \cs_new:Npn \__text_change_case_lower_sigma:nnNw #1#2#3#4 \q__text_recursion_stop
30687 {
30688 \tl_if_head_is_N_type:nTF {#4}
30689 { \__text_change_case_lower_sigma:NnnN #3 }
30690 {
30691 \__text_change_case_store:e
30692 { \char_generate:nn { "03C2 } { \__text_char_catcode:N #3 } }
30693 \__text_change_case_loop:nnw
30694 }
30695 {#1} {#2} #4 \q__text_recursion_stop
30696 }
30697 \cs_new:Npn \__text_change_case_lower_sigma:NnnN #1#2#3#4
30698 {
30699 \__text_change_case_store:e
30700 {
30701 \token_if_letter:NnTF #4
30702 { \char_generate:nn { "03C3 } { \__text_char_catcode:N #1 } }
30703 { \char_generate:nn { "03C2 } { \__text_char_catcode:N #1 } }
30704 }
30705 \__text_change_case_loop:nnw {#2} {#3} #4
30706 }
30707 }

```

In the 8-bit engines, we have to look ahead once we find the first byte of the possible hit.

```

30708 {
30709 \cs_new:Npn \__text_change_case_lower_sigma:nnnN #1#2#3#4
30710 {
30711 \int_compare:nNnTF { '#4 } = { "CE }
30712 { \__text_change_case_lower_sigma:nnnN }
30713 { \__text_change_case_char:nnnN }
30714 {#1} {#2} {#3} #4
30715 }
30716 \cs_new:Npn \__text_change_case_lower_sigma:nnnNN #1#2#3#4#5
30717 {
30718 \int_compare:nNnTF { '#5 } = { "A3 }
30719 { \__text_change_case_lower_sigma:nnw {#2} {#3} }
30720 { \__text_change_case_char:nnnN {#1} {#2} {#3} #4#5 }

```

```

30721     }
30722 \cs_new:Npn \__text_change_case_lower_sigma:nnw #1#2#3 \q__text_recursion_stop
30723 {
30724     \tl_if_head_is_N_type:nTF {#3}
30725     { \__text_change_case_lower_sigma:nnN }
30726     {
30727         \__text_change_case_store:V \c__text_final_sigma_tl
30728         \__text_change_case_loop:nnw
30729     }
30730     {#1} {#2} #3 \q__text_recursion_stop
30731 }
30732 \cs_new:Npn \__text_change_case_lower_sigma:nnN #1#2#3
30733 {
30734     \bool_lazy_or:nnTF
30735     { \token_if_letter_p:N #3 }
30736     {
30737         \bool_lazy_and_p:nn
30738         { \token_if_active_p:N #3 }
30739         { \int_compare_p:nNn { '#3 } > { "80 } }
30740     }
30741     { \__text_change_case_store:V \c__text_sigma_tl }
30742     { \__text_change_case_store:V \c__text_final_sigma_tl }
30743     \__text_change_case_loop:nnw {#1} {#2} #3
30744 }
30745 }

```

For titlecasing, we need to fully expand the new character to see if it is a letter (or active) But that means looking ahead in the 8-bit case, so we have to grab the required tokens up-front. Life is a lot easier for Unicode T_EX's, where we just have one token to worry about. The one wrinkle here is that for look-ahead we'd get into trouble: luckily, only Dutch has that issue.

```

30746 \cs_new:Npx \__text_change_case_char_title:nnN #1#2#3
30747 {
30748     \exp_not:N \bool_if:NTF \l_text_titlecase_check_letter_bool
30749     {
30750         \bool_lazy_or:nnTF
30751         { \sys_if_engine luatex_p: }
30752         { \sys_if_engine xetex_p: }
30753         { \exp_not:N \token_if_letter:NTF #3 }
30754         {
30755             \exp_not:N \bool_lazy_or:nnTF
30756             { \exp_not:N \token_if_letter_p:N #3 }
30757             { \exp_not:N \token_if_active_p:N #3 }
30758         }
30759         { \exp_not:N \use:c { __text_change_case_char_ #1 :nN } }
30760         { \exp_not:N \__text_change_case_char_title:nnnN { title } {#1} }
30761     }
30762     { \exp_not:N \use:c { __text_change_case_char_ #1 :nN } }
30763     {#2} #3
30764 }
30765 \cs_new_eq:NN \__text_change_case_char_titleonly:nnN
30766 \__text_change_case_char_title:nnN
30767 \cs_new:Npn \__text_change_case_char_title:nN #1#2
30768 { \__text_change_case_char_title:nnnN { title } { lower } {#1} #2 }

```

```

30769 \cs_new:Npn \__text_change_case_char_titleonly:nN #1#2
30770 { \__text_change_case_char_title:nnnN { title } { end } {#1} #2 }
30771 \cs_new:Npn \__text_change_case_char_title:nnnN #1#2#3#4
30772 {
30773   \cs_if_exist_use:cF { \__text_change_case_title_ #3 :nnnN }
30774   {
30775     \cs_if_exist_use:cF { \__text_change_case_upper_ #3 :nnnN }
30776     { \__text_change_case_char:nnnN }
30777   }
30778   {#1} {#2} {#3} #4
30779 }

```

For Unicode engines we can handle all characters directly. However, for the 8-bit engines the aim is to deal with (a subset of) Unicode (UTF-8) input. They deal with that by making the upper half of the range active, so we look for that and if found work out how many UTF-8 octets there are to deal with. Those can then be grabbed to reconstruct the full Unicode character, which is then used in a lookup. (As will become obvious below, there is no intention here of covering all of Unicode.)

```

30780 \cs_new:Npn \__text_change_case_char:nnnN #1#2#3#4
30781 {
30782   \token_if_active:NTF #4
30783   { \__text_change_case_store:n {#4} }
30784   {
30785     \__text_change_case_store:e
30786     { \use:c { char_ #1 case :N } #4 }
30787   }
30788   \use:c { \__text_change_case_char_next_ #2 :nn } {#2} {#3}
30789 }
30790 \bool_lazy_or:nnF
30791 { \sys_if_engine luatex_p: }
30792 { \sys_if_engine xetex_p: }
30793 {
30794   \cs_new_eq:NN \__text_change_case_char_aux:nnnN
30795   \__text_change_case_char:nnnN
30796   \cs_gset:Npn \__text_change_case_char:nnnN #1#2#3#4
30797   {
30798     \int_compare:nNnTF { '#4 } > { "80 }
30799     {
30800       \int_compare:nNnTF { '#4 } < { "EO }
30801       { \__text_change_case_char_UTFviii:nnnNN }
30802       {
30803         \int_compare:nNnTF { '#4 } < { "FO }
30804         { \__text_change_case_char_UTFviii:nnnNNN }
30805         { \__text_change_case_char_UTFviii:nnnNNNN }
30806       }
30807       {#1} {#2} {#3} #4
30808     }
30809     { \__text_change_case_char_aux:nnnN {#1} {#2} {#3} #4 }
30810   }
30811   \cs_new:Npn \__text_change_case_char_UTFviii:nnnNN #1#2#3#4#5
30812   { \__text_change_case_char_UTFviii:nnnn {#1} {#2} {#3} {#4#5} }
30813   \cs_new:Npn \__text_change_case_char_UTFviii:nnnNNN #1#2#3#4#5#6
30814   { \__text_change_case_char_UTFviii:nnnn {#1} {#2} {#3} {#4#5#6} }
30815   \cs_new:Npn \__text_change_case_char_UTFviii:nnnNNNN #1#2#3#4#5#6#7

```

```

30816 { \_text_change_case_char_UTFviii:nnnn {#1} {#2} {#3} {#4#5#6#7} }
30817 \cs_new:Npn \_text_change_case_char_UTFviii:nnnn #1#2#3#4
30818 {
30819   \cs_if_exist:cTF { c__text_ #1 case_ \tl_to_str:n {#4} _tl }
30820   {
30821     \_text_change_case_store:v
30822     { c__text_ #1 case_ \tl_to_str:n {#4} _tl }
30823   }
30824   { \_text_change_case_store:n {#4} }
30825   \use:c { __text_change_case_char_next_ #2 :nn } {#2} {#3}
30826 }
30827 }
30828 \cs_new:Npn \_text_change_case_char_next_lower:nn #1#2
30829 { \_text_change_case_loop:nnw {#1} {#2} }
30830 \cs_new_eq:NN \_text_change_case_char_next_upper:nn
30831 \_text_change_case_char_next_lower:nn
30832 \cs_new_eq:NN \_text_change_case_char_next_title:nn
30833 \_text_change_case_char_next_lower:nn
30834 \cs_new_eq:NN \_text_change_case_char_next_titleonly:nn
30835 \_text_change_case_char_next_lower:nn
30836 \cs_new:Npn \_text_change_case_char_next_end:nn #1#2
30837 { \_text_change_case_break:w }

```

(End definition for _text_change_case:nnn and others.)

_text_change_case_upper_de-alt:nnnN
_text_change_case_upper_de-alt:nnnNN

A simple alternative version for German.

```

30838 \bool_lazy_or:nnTF
30839 { \sys_if_engine luatex_p: }
30840 { \sys_if_engine xetex_p: }
30841 {
30842   \cs_new:cpn { __text_change_case_upper_de-alt:nnnN } #1#2#3#4
30843   {
30844     \int_compare:nNnTF { '#4 } = { "00DF }
30845     {
30846       \_text_change_case_store:e
30847       { \char_generate:nn { "1E9E } { \_text_char_catcode:N #4 } }
30848       \use:c { __text_change_case_char_next_ #2 :nn }
30849       {#2} {#3}
30850     }
30851     { \_text_change_case_char:nnnN {#1} {#2} {#3} #4 }
30852   }
30853 }
30854 {
30855   \cs_new:cpx { __text_change_case_upper_de-alt:nnnN } #1#2#3#4
30856   {
30857     \exp_not:N \int_compare:nNnTF { '#4 } = { "00C3 }
30858     {
30859       \exp_not:c { __text_change_case_upper_de-alt:nnnNN }
30860       {#1} {#2} {#3} #4
30861     }
30862     { \exp_not:N \_text_change_case_char:nnnN {#1} {#2} {#3} #4 }
30863   }
30864   \cs_new:cpn { __text_change_case_upper_de-alt:nnnNN } #1#2#3#4#5
30865   {

```

```

30866 \int_compare:nNnTF { '#5 } = { "009F }
30867 {
30868   \__text_change_case_store:V \c__text_grosses_Eszett_tl
30869   \use:c { __text_change_case_char_next_ #2 :nn } {#2} {#3}
30870 }
30871 { \__text_change_case_char:nnnN {#1} {#2} {#3} #4#5 }
30872 }
30873 }

```

(End definition for __text_change_case_upper_de-alt:nnnN and __text_change_case_upper_de-alt:nnnNN.)

```

\__text_change_case_upper_el:nnnN
\__text_change_case_upper_el:nnN
\__text_change_case_upper_el:nnNw
\__text_change_case_upper_el:NnnN
\__text_change_case_upper_el_dialytika:nnN
\__text_change_case_upper_el_dialytika:N
\__text_change_case_upper_el_hiatus:nnNw
\__text_change_case_upper_el_hiatus:nnN
\__text_change_case_upper_el_gobble:nnw
\__text_change_case_upper_el_gobble:nnN
\__text_change_case_if_greek:nTF
\__text_change_case_if_greek_p:n
\__text_change_case_if_greek:nTF
\__text_change_case_if_greek_accent_p:n
\__text_change_case_if_greek_accent:nTF
\__text_change_case_if_greek_diacritic_p:n
\__text_change_case_if_greek_diacritic:nTF
\__text_change_case_if_takes_dialytika:nTF

```

For Greek uppercasing, we need to know if characters *in the Greek range* have accents. That means doing a NFD conversion first, then starting a search. As described by the Unicode CLDR, Greek accents need to be found *after* any U+0308 (diaeresis) and are done in two groups to allow for the canonical ordering. The implementation here follows the data and examples from ICU (<https://sites.google.com/site/icusite/design/case/greek-upper>), although necessarily the implementation is somewhat different.

```

30874 \bool_lazy_or:nnT
30875 { \sys_if_engine luatex_p: }
30876 { \sys_if_engine xetex_p: }
30877 {
30878   \cs_new:Npn \__text_change_case_upper_el:nnnN #1#2#3#4
30879   {
30880     \__text_change_case_if_greek:nTF { '#4 }
30881     {
30882       \exp_args:Ne \__text_change_case_upper_el:nnn
30883       { \char_to_nfd:N #4 } {#2} {#3}
30884     }
30885     { \__text_change_case_char:nnnN {#1} {#2} {#3} #4 }
30886   }
30887   \cs_new:Npn \__text_change_case_upper_el:nnn #1#2#3
30888   { \__text_change_case_upper_el:nnNw {#2} {#3} #1 }

```

At this stage we have the first NFD codepoint as #3. What we need to know is whether after that we have another character token, either from the NFD or directly in the input. If not, we store the changed character at this stage.

```

30889 \cs_new:Npn \__text_change_case_upper_el:nnNw #1#2#3#4 \q__text_recursion_stop
30890 {
30891   \tl_if_head_is_N_type:nTF {#4}
30892   { \__text_change_case_upper_el:NnnN #3 }
30893   {
30894     \__text_change_case_store:e { \char_uppercase:N #3 }
30895     \__text_change_case_loop:nnw
30896   }
30897   {#1} {#2} #4 \q__text_recursion_stop
30898 }

```

Now, we check the detail of the next codepoint: again we filter out the not-a-char cases, before checking if it's an dialytika, accent or diacritic. (The latter do not have the same hiatus behavior as accents.)

```

30899 \cs_new:Npn \__text_change_case_upper_el:NnnN #1#2#3#4
30900 {
30901   \token_if_cs:NTF #4
30902   {

```

```

30903         \_text_change_case_store:e { \char_uppercase:N #1 }
30904         \_text_change_case_loop:nnw {#2} {#3} #4
30905     }
30906     {
30907         \int_compare:nNnTF { '#4 } = { "0308 }
30908         { \_text_change_case_upper_el_dialytika:nnN {#2} {#3} #1 }
30909         {
30910             \_text_change_case_if_greek_accent:nTF { '#4 }
30911             { \_text_change_case_upper_el_hiatus:nnNw {#2} {#3} #1 }
30912             {
30913                 \_text_change_case_if_greek_diacritic:nTF { '#4 }
30914                 {
30915                     \_text_change_case_store:e { \char_uppercase:N #1 }
30916                     \_text_change_case_loop:nnw {#2} {#3}
30917                 }
30918                 {
30919                     \_text_change_case_store:e { \char_uppercase:N #1 }
30920                     \_text_change_case_loop:nnw {#2} {#3} #4
30921                 }
30922             }
30923         }
30924     }
30925 }

```

We handle *dialytika* in parts as it's also needed for the hiatus. We know only two letters take it, so we can shortcut here on the second part of the tests.

```

30926 \cs_new:Npn \_text_change_case_upper_el_dialytika:nnN #1#2#3
30927 {
30928     \_text_change_case_if_takes_dialytika:nTF { '#3 }
30929     { \_text_change_case_upper_el_dialytika:N #3 }
30930     { \_text_change_case_store:e { \char_uppercase:N #3 } }
30931     \_text_change_case_upper_el_gobble:nnw {#1} {#2}
30932 }
30933 \cs_new:Npn \_text_change_case_upper_el_dialytika:N #1
30934 {
30935     \_text_change_case_store:e
30936     {
30937         \bool_lazy_or:nnTF
30938         { \int_compare_p:nNn { '#1 } = { "0399 } }
30939         { \int_compare_p:nNn { '#1 } = { "03B9 } }
30940         { \char_generate:nn { "03AA } { \_text_char_catcode:N #1 } }
30941         { \char_generate:nn { "03AB } { \_text_char_catcode:N #1 } }
30942     }
30943 }

```

Adding a hiatus needs some of the same ideas, but if there is not one we skip this code point, hence needing a separate function.

```

30944 \cs_new:Npn \_text_change_case_upper_el_hiatus:nnNw
30945 #1#2#3#4 \q__text_recursion_stop
30946 {
30947     \_text_change_case_store:e { \char_uppercase:N #3 }
30948     \tl_if_head_is_N_type:nTF {#4}
30949     { \_text_change_case_upper_el_hiatus:nnN }
30950     { \_text_change_case_loop:nnw }
30951     {#1} {#2} #4 \q__text_recursion_stop

```

```

30952     }
30953 \cs_new:Npn \__text_change_case_upper_el_hiatus:nnN #1#2#3
30954 {
30955     \token_if_cs:NTF #3
30956     { \__text_change_case_loop:nnw {#1} {#2} #3 }
30957     {
30958         \__text_change_case_if_takes_dialytika:nTF { '#3 }
30959         {
30960             \__text_change_case_upper_el_dialytika:N #3
30961             \__text_change_case_upper_el_gobble:nnw {#1} {#2}
30962         }
30963         { \__text_change_case_loop:nnw {#1} {#2} #3 }
30964     }
30965 }

```

For clearing out trailing combining marks after we have dealt with the first one.

```

30966 \cs_new:Npn \__text_change_case_upper_el_gobble:nnw
30967 #1#2#3 \q__text_recursion_stop
30968 {
30969     \tl_if_head_is_N_type:nTF {#3}
30970     { \__text_change_case_upper_el_gobble:nnN }
30971     { \__text_change_case_loop:nnw }
30972     {#1} {#2} #3 \q__text_recursion_stop
30973 }
30974 \cs_new:Npn \__text_change_case_upper_el_gobble:nnN #1#2#3
30975 {
30976     \bool_lazy_or:nnTF
30977     { \token_if_cs_p:N #3 }
30978     {
30979         ! \bool_lazy_or_p:nn
30980         { \__text_change_case_if_greek_accent_p:n { '#3 } }
30981         { \__text_change_case_if_greek_diacritic_p:n { '#3 } }
30982     }
30983     { \__text_change_case_loop:nnw {#1} {#2} #3 }
30984     { \__text_change_case_upper_el_gobble:nnw {#1} {#2} }
30985 }
30986 }

```

Luckily the Greek range is limited and clear.

```

30987 \prg_new_conditional:Npnn \__text_change_case_if_greek:n #1 { TF }
30988 {
30989     \if_int_compare:w #1 < "0370 \exp_stop_f:
30990     \prg_return_false:
30991 }else:
30992     \if_int_compare:w #1 > "03FF \exp_stop_f:
30993     \if_int_compare:w #1 < "1F00 \exp_stop_f:
30994     \prg_return_false:
30995 }else:
30996     \if_int_compare:w #1 > "1FFF \exp_stop_f:
30997     \prg_return_false:
30998 }else:
30999     \prg_return_true:
31000 \fi:
31001 \fi:
31002 }else:

```

```

31003         \prg_return_true:
31004     \fi:
31005 \fi:
31006 }

```

We follow ICU in adding a few extras to the accent list here.

```

31007 \prg_new_conditional:Npnn \_text_change_case_if_greek_accent:n #1 { TF , p }
31008 {
31009     \if_int_compare:w #1 = "0300 \exp_stop_f:
31010         \prg_return_true:
31011     \else:
31012         \if_int_compare:w #1 = "0301 \exp_stop_f:
31013         \prg_return_true:
31014     \else:
31015         \if_int_compare:w #1 = "0342 \exp_stop_f:
31016         \prg_return_true:
31017     \else:
31018         \if_int_compare:w #1 = "0302 \exp_stop_f:
31019         \prg_return_true:
31020     \else:
31021         \if_int_compare:w #1 = "0303 \exp_stop_f:
31022         \prg_return_true:
31023     \else:
31024         \if_int_compare:w #1 = "0311 \exp_stop_f:
31025         \prg_return_true:
31026     \else:
31027         \prg_return_false:
31028     \fi:
31029 \fi:
31030 \fi:
31031 \fi:
31032 \fi:
31033 \fi:
31034 }
31035 \prg_new_conditional:Npnn \_text_change_case_if_greek_diacritic:n
31036 #1 { TF , p }
31037 {
31038     \if_int_compare:w #1 = "0304 \exp_stop_f:
31039         \prg_return_true:
31040     \else:
31041         \if_int_compare:w #1 = "0306 \exp_stop_f:
31042         \prg_return_true:
31043     \else:
31044         \if_int_compare:w #1 = "0313 \exp_stop_f:
31045         \prg_return_true:
31046     \else:
31047         \if_int_compare:w #1 = "0314 \exp_stop_f:
31048         \prg_return_true:
31049     \else:
31050         \if_int_compare:w #1 = "0343 \exp_stop_f:
31051         \prg_return_true:
31052     \else:
31053         \prg_return_false:
31054     \fi:
31055 \fi:

```



```

31056         \fi:
31057     \fi:
31058 \fi:
31059 }
31060 \prg_new_conditional:Npnn \_text_change_case_if_takes_dialytika:n #1 { TF }
31061 {
31062     \if_int_compare:w #1 = "0399 \exp_stop_f:
31063     \prg_return_true:
31064 \else:
31065     \if_int_compare:w #1 = "03B9 \exp_stop_f:
31066     \prg_return_true:
31067 \else:
31068     \if_int_compare:w #1 = "03A5 \exp_stop_f:
31069     \prg_return_true:
31070 \else:
31071     \if_int_compare:w #1 = "03C5 \exp_stop_f:
31072     \prg_return_true:
31073 \else:
31074     \prg_return_false:
31075 \fi:
31076 \fi:
31077 \fi:
31078 \fi:
31079 }

```

(End definition for _text_change_case_upper_el:nnnN and others.)

_text_change_case_boundary_upper_el:Nnnw There is one special case in Greek that needs to be picked up based on being an isolated letter. We do that using a test similar to final sigma, but it has to fire off from the space grabber.

_text_change_case_boundary_upper_el:nnN

_text_change_case_boundary_upper_el:nnw

_text_change_case_boundary_upper_el:NnnN

```

31080 \bool_lazy_or:nnT
31081 { \sys_if_engine luatex_p: }
31082 { \sys_if_engine xetex_p: }
31083 {
31084     \cs_new:Npn \_text_change_case_boundary_upper_el:Nnnw
31085     #1#2#3#4 \q_text_recursion_stop
31086     {
31087         \tl_if_head_is_N_type:nTF {#4}
31088         { \_text_change_case_boundary_upper_el:nnN }
31089         { \_text_change_case_loop:nnw }
31090         {#2} {#3} #4 \q_text_recursion_stop
31091     }
31092 \cs_new:Npn \_text_change_case_boundary_upper_el:nnN #1#2#3
31093 {
31094     \bool_lazy_or:nnTF
31095     { \token_if_cs_p:N #3 }
31096     {
31097         ! \bool_lazy_or_p:nn
31098         { \int_compare_p:nNn { '#3 } = { "03AE } }
31099         { \int_compare_p:nNn { '#3 } = { "1F22 } }
31100     }
31101     { \_text_change_case_loop:nnw }
31102     { \_text_change_case_boundary_upper_el:nnw }
31103     {#1} {#2} #3

```

```

31104     }
31105 \cs_new:Npn \__text_change_case_boundary_upper_el:nnNw
31106   #1#2#3#4 \q__text_recursion_stop
31107   {
31108     \tl_if_head_is_N_type:nTF {#4}
31109     { \__text_change_case_boundary_upper_el:NnnN #3 }
31110     {
31111       \__text_change_case_store:e
31112       { \char_generate:nn { "0389 } { \__text_char_catcode:N #3 } }
31113       \__text_change_case_loop:nnw
31114     }
31115     {#1} {#2} #4 \q__text_recursion_stop
31116   }
31117 \cs_new:Npn \__text_change_case_boundary_upper_el:NnnN #1#2#3#4
31118   {
31119     \token_if_letter:NTF #4
31120     { \__text_change_case_loop:nnw {#2} {#3} #1#4 }
31121     {
31122       \__text_change_case_store:e
31123       { \char_generate:nn { "0389 } { \__text_char_catcode:N #1 } }
31124       \__text_change_case_loop:nnw {#2} {#3} #4
31125     }
31126   }
31127 }

```

(End definition for __text_change_case_boundary_upper_el:Nnnw and others.)

__text_change_case_title_el:nnnN Titlecasing retains accents, but to prevent the uppercasing code from kicking in, there has to be an explicit function here.

```

31128 \bool_lazy_or:nnT
31129 { \sys_if_engine luatex_p: }
31130 { \sys_if_engine xetex_p: }
31131 {
31132   \cs_new:Npn \__text_change_case_title_el:nnnN #1#2#3#4
31133   { \__text_change_case_char:nnnN {#1} {#2} {#3} #4 }
31134 }

```

(End definition for __text_change_case_title_el:nnnN.)

__text_change_cases_lower_lt:nnnN For Lithuanian, the issue to be dealt with is dots over lower case letters: these should be present if there is another accent. The first step is a simple match attempt: look for the three uppercase accented letters which should gain a dot-above char in their lowercase form.

```

\__text_change_cases_lower_lt_auxi:nnnN
\__text_change_cases_lower_lt_auxii:nnnN
\__text_change_case_lower_lt:nnw
\__text_change_case_lower_lt:nnN
31135 \bool_lazy_or:nnT
31136 { \sys_if_engine luatex_p: }
31137 { \sys_if_engine xetex_p: }
31138 {
31139   \cs_new:Npn \__text_change_case_lower_lt:nnnN #1#2#3#4
31140   {
31141     \exp_args:Ne \__text_change_case_lower_lt_auxi:nnnN
31142     {
31143       \int_case:nn { '#4 }
31144       {
31145         { "00CC } { "0300 }

```

```

31146         { "00CD } { "0301 }
31147         { "0128 } { "0303 }
31148     }
31149 }
31150     {#2} {#3} #4
31151 }

```

If there was a hit, output the result with the dot-above and move on. Otherwise, look for one of the three letters that can take a combining accent: I, J and I-ogonek.

```

31152 \cs_new:Npn \__text_change_case_lower_lt_auxi:nnnN #1#2#3#4
31153 {
31154     \tl_if_blank:nTF {#1}
31155     {
31156         \exp_args:Ne \__text_change_case_lower_lt_auxii:nnnN
31157         {
31158             \int_case:nn { '#4 }
31159             {
31160                 { "0049 } { "0069 }
31161                 { "004A } { "006A }
31162                 { "012E } { "012F }
31163             }
31164         }
31165         {#2} {#3} #4
31166     }
31167     {
31168         \__text_change_case_store:e
31169         {
31170             \char_generate:nn { "0069 } { \__text_char_catcode:N #4 }
31171             \char_generate:nn { "0307 } { \__text_char_catcode:N #4 }
31172             \char_generate:nn {#1} { \__text_char_catcode:N #4 }
31173         }
31174         \__text_change_case_loop:nnw {#2} {#3}
31175     }
31176 }

```

Again, branch depending on a hit. If there is one, we output the character then need to look for a combining accent: as usual, we need to be aware of the loop situation.

```

31177 \cs_new:Npn \__text_change_case_lower_lt_auxii:nnnN #1#2#3#4
31178 {
31179     \tl_if_blank:nTF {#1}
31180     { \__text_change_case_lower_sigma:nnnN {#2} {#2} {#3} #4 }
31181     {
31182         \__text_change_case_store:e
31183         { \char_generate:nn {#1} { \__text_char_catcode:N #4 } }
31184         \__text_change_case_lower_lt:nnw {#2} {#3}
31185     }
31186 }
31187 \cs_new:Npn \__text_change_case_lower_lt:nnw #1#2#3 \q__text_recursion_stop
31188 {
31189     \tl_if_head_is_N_type:nTF {#3}
31190     { \__text_change_case_lower_lt:nnN }
31191     { \__text_change_case_loop:nnw }
31192     {#1} {#2} #3 \q__text_recursion_stop
31193 }
31194 \cs_new:Npn \__text_change_case_lower_lt:nnN #1#2#3

```

```

31195 {
31196   \bool_lazy_and:nnT
31197   { ! \token_if_cs_p:N #3 }
31198   {
31199     \bool_lazy_any_p:n
31200     {
31201       { \int_compare_p:nNn { '#3 } = { "0300 } }
31202       { \int_compare_p:nNn { '#3 } = { "0301 } }
31203       { \int_compare_p:nNn { '#3 } = { "0303 } }
31204     }
31205   }
31206   {
31207     \__text_change_case_store:e
31208     { \char_generate:nn { "0307 } { \__text_char_catcode:N #3 } }
31209   }
31210   \__text_change_case_loop:nnw {#1} {#2} #3
31211 }
31212 }

```

(End definition for __text_change_cases_lower_lt:nnnN and others.)

__text_change_cases_upper_lt:nnnN
 __text_change_cases_upper_lt_aux:nnnN
 __text_change_case_upper_lt:nnw
 __text_change_case_upper_lt:nnN

The uppercasing version: first find i/j/i-ogonek, then look for the combining char: drop it if present.

```

31213 \bool_lazy_or:nnT
31214 { \sys_if_engine luatex_p: }
31215 { \sys_if_engine xetex_p: }
31216 {
31217   \cs_new:Npn \__text_change_case_upper_lt:nnnN #1#2#3#4
31218   {
31219     \exp_args:Ne \__text_change_case_upper_lt_aux:nnnN
31220     {
31221       \int_case:nn { '#4 }
31222       {
31223         { "0069 } { "0049 }
31224         { "006A } { "004A }
31225         { "012F } { "012E }
31226       }
31227     }
31228     {#2} {#3} #4
31229   }
31230   \cs_new:Npn \__text_change_case_upper_lt_aux:nnnN #1#2#3#4
31231   {
31232     \tl_if_blank:nTF {#1}
31233     { \__text_change_case_char:nnnN { upper } {#2} {#3} #4 }
31234     {
31235       \__text_change_case_store:e
31236       { \char_generate:nn {#1} { \__text_char_catcode:N #4 } }
31237       \__text_change_case_upper_lt:nnw {#2} {#3}
31238     }
31239   }
31240   \cs_new:Npn \__text_change_case_upper_lt:nnw #1#2#3 \q__text_recursion_stop
31241   {
31242     \tl_if_head_is_N_type:nTF {#3}
31243     { \__text_change_case_upper_lt:nnN }

```

```

31244         { \use:c { __text_change_case_char_next_ #1 :nn } }
31245         {#1} {#2} #3 \q__text_recursion_stop
31246     }
31247 \cs_new:Npn \__text_change_case_upper_lt:nnN #1#2#3
31248 {
31249     \bool_lazy_and:nnTF
31250     { ! \token_if_cs_p:N #3 }
31251     { \int_compare_p:nNn { '#3 } = { "0307 } }
31252     { \use:c { __text_change_case_char_next_ #1 :nn } {#1} {#2} }
31253     { \use:c { __text_change_case_char_next_ #1 :nn } {#1} {#2} #3 }
31254 }
31255 }

```

(End definition for __text_change_cases_upper_lt:nnnN and others.)

__text_change_case_title_nl:nnnN For Dutch, there is a single look-ahead test for ij when title casing. If the appropriate letters are found, produce IJ and gobble the j/J.

```

\__text_change_case_title_nl:nnw
\__text_change_case_title_nl:nnN
31256 \cs_new:Npn \__text_change_case_title_nl:nnnN #1#2#3#4
31257 {
31258     \bool_lazy_or:nnTF
31259     { \int_compare_p:nNn { '#4 } = { "0049 } }
31260     { \int_compare_p:nNn { '#4 } = { "0069 } }
31261     {
31262         \__text_change_case_store:e
31263         { \char_generate:nn { "0049 } { \__text_char_catcode:N #4 } }
31264         \__text_change_case_title_nl:nnw {#2} {#3}
31265     }
31266     { \__text_change_case_char:nnnN {#1} {#2} {#3} #4 }
31267 }
31268 \cs_new:Npn \__text_change_case_title_nl:nnw #1#2#3 \q__text_recursion_stop
31269 {
31270     \tl_if_head_is_N_type:nTF {#3}
31271     { \__text_change_case_title_nl:nnN }
31272     { \use:c { __text_change_case_char_next_ #1 :nn } }
31273     {#1} {#2} #3 \q__text_recursion_stop
31274 }
31275 \cs_new:Npn \__text_change_case_title_nl:nnN #1#2#3
31276 {
31277     \bool_lazy_and:nnTF
31278     { ! \token_if_cs_p:N #3 }
31279     {
31280         \bool_lazy_or_p:nn
31281         { \int_compare_p:nNn { '#3 } = { "004A } }
31282         { \int_compare_p:nNn { '#3 } = { "006A } }
31283     }
31284     {
31285         \__text_change_case_store:e
31286         { \char_generate:nn { "004A } { \__text_char_catcode:N #3 } }
31287         \use:c { __text_change_case_char_next_ #1 :nn } {#1} {#2}
31288     }
31289     { \use:c { __text_change_case_char_next_ #1 :nn } {#1} {#2} #3 }
31290 }

```

(End definition for __text_change_case_title_nl:nnnN, __text_change_case_title_nl:nnw, and __text_change_case_title_nl:nnN.)

```

\__text_change_case_lower_tr:nnnN
\__text_change_case_lower_tr:nnNw
\__text_change_case_lower_tr:NnnN
\__text_change_case_lower_tr:nnnNN

```

The Turkic languages need special treatment for dotted-i and dotless-i. The lower casing rule can be expressed in terms of searching first for either a dotless-I or a dotted-I. In the latter case the mapping is easy, but in the former there is a second stage search.

```

31291 \bool_lazy_or:nnTF
31292 { \sys_if_engine_luatex_p: }
31293 { \sys_if_engine_xetex_p: }
31294 {
31295   \cs_new:Npn \__text_change_case_lower_tr:nnnN #1#2#3#4
31296   {
31297     \int_compare:nNnTF { '#4 } = { "0049 }
31298     { \__text_change_case_lower_tr:nnNw {#1} {#3} #4 }
31299     {
31300       \int_compare:nNnTF { '#4 } = { "0130 }
31301       {
31302         \__text_change_case_store:e
31303         { \char_generate:nn { "0069 } { \__text_char_catcode:N #4 } }
31304         \__text_change_case_loop:nnw {#1} {#3}
31305       }
31306       { \__text_change_case_lower_sigma:nnnN {#1} {#2} {#3} #4 }
31307     }
31308   }

```

After a dotless-I there may be a dot-above character. If there is then a dotted-i should be produced, otherwise output a dotless-i. When the combination is found both the dotless-I and the dot-above char have to be removed from the input.

```

31309 \cs_new:Npn \__text_change_case_lower_tr:nnNw #1#2#3#4 \q__text_recursion_stop
31310 {
31311   \tl_if_head_is_N_type:nTF {#4}
31312   { \__text_change_case_lower_tr:NnnN #3 }
31313   {
31314     \__text_change_case_store:e
31315     { \char_generate:nn { "0131 } { \__text_char_catcode:N #3 } }
31316     \__text_change_case_loop:nnw
31317     {#1} {#2} #4 \q__text_recursion_stop
31318   }
31319 }
31320 \cs_new:Npn \__text_change_case_lower_tr:NnnN #1#2#3#4
31321 {
31322   \bool_lazy_or:nnTF
31323   { \token_if_cs_p:N #4 }
31324   { ! \int_compare_p:nNn { '#4 } = { "0307 } }
31325   {
31326     \__text_change_case_store:e
31327     { \char_generate:nn { "0131 } { \__text_char_catcode:N #1 } }
31328     \__text_change_case_loop:nnw {#2} {#3} #4
31329   }
31330   {
31331     \__text_change_case_store:e
31332     { \char_generate:nn { "0069 } { \__text_char_catcode:N #1 } }
31333     \__text_change_case_loop:nnw {#2} {#3}
31334   }
31335 }
31336 }

```

For 8-bit engines, dot-above is not available so there is a simple test for an upper-case I. Then we can look for the UTF-8 representation of an upper case dotted-I without the combining char. If it's not there, preserve the UTF-8 sequence as-is. With 8bit engines, we cannot completely preserve category codes, so we have to make some assumptions: output a “normal” i for the dotted case. As the original character here is catcode-13, we have to make a choice about handling of i: generate a “normal” one.

```

31337 {
31338   \cs_new:Npn \__text_change_case_lower_tr:nnnN #1#2#3#4
31339   {
31340     \int_compare:nNnTF { '#4 } = { "0049 }
31341     {
31342       \__text_change_case_store:V \c__text_dotless_i_tl
31343       \__text_change_case_loop:nnw {#1} {#3}
31344     }
31345     {
31346       \int_compare:nNnTF { '#4 } = { "00C4 }
31347       { \__text_change_case_lower_tr:nnnNN {#1} {#2} {#3} #4 }
31348       { \__text_change_case_char:nnnN {#1} {#2} {#3} #4 }
31349     }
31350   }
31351   \cs_new:Npn \__text_change_case_lower_tr:nnnNN #1#2#3#4#5
31352   {
31353     \int_compare:nNnTF { '#5 } = { "00B0 }
31354     {
31355       \__text_change_case_store:e
31356       {
31357         \char_generate:nn { "0069 }
31358         { \char_value_catcode:n { "0069 } }
31359       }
31360       \__text_change_case_loop:nnw {#1} {#3}
31361     }
31362     { \__text_change_case_char:nnnN {#1} {#2} {#3} #4#5 }
31363   }
31364 }

```

(End definition for __text_change_case_lower_tr:nnnN and others.)

__text_change_case_upper_tr:nnnN

Uppercasing is easier: just one exception with no context.

```

31365 \cs_new:Npx \__text_change_case_upper_tr:nnnN #1#2#3#4
31366 {
31367   \exp_not:N \int_compare:nNnTF { '#4 } = { "0069 }
31368   {
31369     \bool_lazy_or:nnTF
31370     { \sys_if_engine luatex_p: }
31371     { \sys_if_engine xetex_p: }
31372     {
31373       \exp_not:N \__text_change_case_store:e
31374       {
31375         \exp_not:N \char_generate:nn { "0130 }
31376         { \exp_not:N \__text_char_catcode:N #4 }
31377       }
31378     }
31379     {
31380       \exp_not:N \__text_change_case_store:V

```

```

31381         \exp_not:N \c__text_dotted_I_tl
31382     }
31383     \exp_not:N \use:c { __text_change_case_char_next_ #2 :nn } {#2} {#3}
31384 }
31385 { \exp_not:N \__text_change_case_char:nnnN {#1} {#2} {#3} #4 }
31386 }

```

(End definition for __text_change_case_upper_tr:nnnN.)

__text_change_case_lower_az:nnnN

Straight copies.

__text_change_case_upper_az:nnnN

```

31387 \cs_new_eq:NN \__text_change_case_lower_az:nnnN
31388 \__text_change_case_lower_tr:nnnN
31389 \cs_new_eq:NN \__text_change_case_upper_az:nnnN
31390 \__text_change_case_upper_tr:nnnN

```

(End definition for __text_change_case_lower_az:nnnN and __text_change_case_upper_az:nnnN.)

48.2 Case changing data for 8-bit engines

For cases where there is an 8-bit option in the T1 font set up, a variant is provided in both cases. There are also a few extras for LGR.

```

\c__text_dotless_i_tl
\c__text_dotted_I_tl
\c__text_i_ogonek_tl
\c__text_I_ogonek_tl
\c__text_final_sigma_tl
\c__text_sigma_tl
\c__text_grosses_Eszett_tl
31391 \group_begin:
31392   \bool_lazy_or:nnF
31393   { \sys_if_engine_luatex_p: }
31394   { \sys_if_engine_xetex_p: }
31395   {
31396     \cs_set_protected:Npn \__text_tmp:w #1#2
31397     {
31398       \group_begin:
31399       \cs_set_protected:Npn \__text_tmp:w ##1##2##3##4
31400       {
31401         \tl_const:Nx #1
31402         {
31403           \exp_after:wN \exp_after:wN \exp_after:wN
31404           \exp_not:N \char_generate:nn {##1} { 13 }
31405           \exp_after:wN \exp_after:wN \exp_after:wN
31406           \exp_not:N \char_generate:nn {##2} { 13 }
31407           \tl_if_blank:nF {##3}
31408           {
31409             \exp_after:wN \exp_after:wN \exp_after:wN
31410             \exp_not:N \char_generate:nn {##3} { 13 }
31411           }
31412         }
31413       }
31414       \use:x
31415       { \__text_tmp:w \char_to_utfviii_bytes:n { "#2 } }
31416     }
31417   }
31418   \__text_tmp:w \c__text_dotless_i_tl      { 0131 }
31419   \__text_tmp:w \c__text_dotted_I_tl      { 0130 }
31420   \__text_tmp:w \c__text_i_ogonek_tl      { 012F }
31421   \__text_tmp:w \c__text_I_ogonek_tl      { 012E }
31422   \__text_tmp:w \c__text_final_sigma_tl   { 03C2 }
31423   \__text_tmp:w \c__text_sigma_tl         { 03C3 }

```



```

31424     \__text_tmp:w \c__text_grosses_Eszett_tl { 1E9E }
31425   }
31426 \group_end:

```

(End definition for `\c__text_dotless_i_tl` and others.)

For 8-bit engines we now need to define the case-change data for the multi-octet mappings. This data is here not in the `char` module as the multi-byte nature means they are never N-type. These need a list of what code points are doable in T1 so the list is hard coded (there's no saving in loading the mappings dynamically). All of the straight-forward ones have two octets, so that is taken as read.

```

31427 \group_begin:
31428   \bool_lazy_or:nnF
31429   { \sys_if_engine luatex_p: }
31430   { \sys_if_engine xetex_p: }
31431   {
31432     \cs_set_protected:Npn \__text_loop:nn #1#2
31433     {
31434       \quark_if_recursion_tail_stop:n {#1}
31435       \use:x
31436       {
31437         \__text_tmp:w
31438         \char_to_utfviii_bytes:n { "#1 }
31439         \char_to_utfviii_bytes:n { "#2 }
31440       }
31441       \__text_loop:nn
31442     }
31443     \cs_set_protected:Npn \__text_tmp:nnnn #1#2#3#4#5
31444     {
31445       \tl_const:cx
31446       {
31447         c__text_ #1 case_
31448         \char_generate:nn {#2} { 12 }
31449         \char_generate:nn {#3} { 12 }
31450         _tl
31451       }
31452       {
31453         \exp_after:wN \exp_after:wN \exp_after:wN
31454         \exp_not:N \char_generate:nn {#4} { 13 }
31455         \exp_after:wN \exp_after:wN \exp_after:wN
31456         \exp_not:N \char_generate:nn {#5} { 13 }
31457       }
31458     }
31459     \cs_set_protected:Npn \__text_tmp:w #1#2#3#4#5#6#7#8
31460     {
31461       \tl_const:cx
31462       {
31463         c__text_lowercase_
31464         \char_generate:nn {#1} { 12 }
31465         \char_generate:nn {#2} { 12 }
31466         _tl
31467       }
31468       {
31469         \exp_after:wN \exp_after:wN \exp_after:wN
31470         \exp_not:N \char_generate:nn {#5} { 13 }

```

```

31471         \exp_after:wN \exp_after:wN \exp_after:wN
31472         \exp_not:N \char_generate:nn {#6} { 13 }
31473     }
31474     \__text_tmp:nnnn { upper } {#5} {#6} {#1} {#2}
31475     \__text_tmp:nnnn { title } {#5} {#6} {#1} {#2}
31476 }
31477 \__text_loop:nn
31478 { 00C0 } { 00E0 }
31479 { 00C1 } { 00E1 }
31480 { 00C2 } { 00E2 }
31481 { 00C3 } { 00E3 }
31482 { 00C4 } { 00E4 }
31483 { 00C5 } { 00E5 }
31484 { 00C6 } { 00E6 }
31485 { 00C7 } { 00E7 }
31486 { 00C8 } { 00E8 }
31487 { 00C9 } { 00E9 }
31488 { 00CA } { 00EA }
31489 { 00CB } { 00EB }
31490 { 00CC } { 00EC }
31491 { 00CD } { 00ED }
31492 { 00CE } { 00EE }
31493 { 00CF } { 00EF }
31494 { 00D0 } { 00F0 }
31495 { 00D1 } { 00F1 }
31496 { 00D2 } { 00F2 }
31497 { 00D3 } { 00F3 }
31498 { 00D4 } { 00F4 }
31499 { 00D5 } { 00F5 }
31500 { 00D6 } { 00F6 }
31501 { 00D8 } { 00F8 }
31502 { 00D9 } { 00F9 }
31503 { 00DA } { 00FA }
31504 { 00DB } { 00FB }
31505 { 00DC } { 00FC }
31506 { 00DD } { 00FD }
31507 { 00DE } { 00FE }
31508 { 0100 } { 0101 }
31509 { 0102 } { 0103 }
31510 { 0104 } { 0105 }
31511 { 0106 } { 0107 }
31512 { 0108 } { 0109 }
31513 { 010A } { 010B }
31514 { 010C } { 010D }
31515 { 010E } { 010F }
31516 { 0110 } { 0111 }
31517 { 0112 } { 0113 }
31518 { 0114 } { 0115 }
31519 { 0116 } { 0117 }
31520 { 0118 } { 0119 }
31521 { 011A } { 011B }
31522 { 011C } { 011D }
31523 { 011E } { 011F }
31524 { 0120 } { 0121 }

```

31525	{ 0122 }	{ 0123 }
31526	{ 0124 }	{ 0125 }
31527	{ 0128 }	{ 0129 }
31528	{ 012A }	{ 012B }
31529	{ 012C }	{ 012D }
31530	{ 012E }	{ 012F }
31531	{ 0132 }	{ 0133 }
31532	{ 0134 }	{ 0135 }
31533	{ 0136 }	{ 0137 }
31534	{ 0139 }	{ 013A }
31535	{ 013B }	{ 013C }
31536	{ 013E }	{ 013F }
31537	{ 0141 }	{ 0142 }
31538	{ 0143 }	{ 0144 }
31539	{ 0145 }	{ 0146 }
31540	{ 0147 }	{ 0148 }
31541	{ 014A }	{ 014B }
31542	{ 014C }	{ 014D }
31543	{ 014E }	{ 014F }
31544	{ 0150 }	{ 0151 }
31545	{ 0152 }	{ 0153 }
31546	{ 0154 }	{ 0155 }
31547	{ 0156 }	{ 0157 }
31548	{ 0158 }	{ 0159 }
31549	{ 015A }	{ 015B }
31550	{ 015C }	{ 015D }
31551	{ 015E }	{ 015F }
31552	{ 0160 }	{ 0161 }
31553	{ 0162 }	{ 0163 }
31554	{ 0164 }	{ 0165 }
31555	{ 0168 }	{ 0169 }
31556	{ 016A }	{ 016B }
31557	{ 016C }	{ 016D }
31558	{ 016E }	{ 016F }
31559	{ 0170 }	{ 0171 }
31560	{ 0172 }	{ 0173 }
31561	{ 0174 }	{ 0175 }
31562	{ 0176 }	{ 0177 }
31563	{ 0178 }	{ 00FF }
31564	{ 0179 }	{ 017A }
31565	{ 017B }	{ 017C }
31566	{ 017D }	{ 017E }
31567	{ 01CD }	{ 01CE }
31568	{ 01CF }	{ 01D0 }
31569	{ 01D1 }	{ 01D2 }
31570	{ 01D3 }	{ 01D4 }
31571	{ 01E2 }	{ 01E3 }
31572	{ 01E6 }	{ 01E7 }
31573	{ 01E8 }	{ 01E9 }
31574	{ 01EA }	{ 01EB }
31575	{ 01F4 }	{ 01F5 }
31576	{ 0218 }	{ 0219 }
31577	{ 021A }	{ 021B }

Add T2 (Cyrillic) as this is doable using a classical \MakeUppercase approach.

31578	{ 0400 }	{ 0450 }
31579	{ 0401 }	{ 0451 }
31580	{ 0402 }	{ 0452 }
31581	{ 0403 }	{ 0453 }
31582	{ 0404 }	{ 0454 }
31583	{ 0405 }	{ 0455 }
31584	{ 0406 }	{ 0456 }
31585	{ 0407 }	{ 0457 }
31586	{ 0408 }	{ 0458 }
31587	{ 0409 }	{ 0459 }
31588	{ 040A }	{ 045A }
31589	{ 040B }	{ 045B }
31590	{ 040C }	{ 045C }
31591	{ 040D }	{ 045D }
31592	{ 040E }	{ 045E }
31593	{ 040F }	{ 045F }
31594	{ 0410 }	{ 0430 }
31595	{ 0411 }	{ 0431 }
31596	{ 0412 }	{ 0432 }
31597	{ 0413 }	{ 0433 }
31598	{ 0414 }	{ 0434 }
31599	{ 0415 }	{ 0435 }
31600	{ 0416 }	{ 0436 }
31601	{ 0417 }	{ 0437 }
31602	{ 0418 }	{ 0438 }
31603	{ 0419 }	{ 0439 }
31604	{ 041A }	{ 043A }
31605	{ 041B }	{ 043B }
31606	{ 041C }	{ 043C }
31607	{ 041D }	{ 043D }
31608	{ 041E }	{ 043E }
31609	{ 041F }	{ 043F }
31610	{ 0420 }	{ 0440 }
31611	{ 0421 }	{ 0441 }
31612	{ 0422 }	{ 0442 }
31613	{ 0423 }	{ 0443 }
31614	{ 0424 }	{ 0444 }
31615	{ 0425 }	{ 0445 }
31616	{ 0426 }	{ 0446 }
31617	{ 0427 }	{ 0447 }
31618	{ 0428 }	{ 0448 }
31619	{ 0429 }	{ 0449 }
31620	{ 042A }	{ 044A }
31621	{ 042B }	{ 044B }
31622	{ 042C }	{ 044C }
31623	{ 042D }	{ 044D }
31624	{ 042E }	{ 044E }
31625	{ 042F }	{ 044F }

Greek support: everything in the two-octet range.

31626	{ 0370 }	{ 0371 }
31627	{ 0372 }	{ 0373 }
31628	{ 0376 }	{ 0377 }

```

31629      { 03FD } { 037B }
31630      { 03FE } { 037C }
31631      { 03FF } { 037D }
31632      { 0386 } { 03AC }
31633      { 0388 } { 03AD }
31634      { 0389 } { 03AE }
31635      { 038A } { 03AF }
31636      { 0391 } { 03B1 }
31637      { 0392 } { 03B2 }
31638      { 0393 } { 03B3 }
31639      { 0394 } { 03B4 }
31640      { 0395 } { 03B5 }
31641      { 0396 } { 03B6 }
31642      { 0397 } { 03B7 }
31643      { 0398 } { 03B8 }
31644      { 0399 } { 03B9 }
31645      { 039A } { 03BA }
31646      { 039B } { 03BB }
31647      { 039C } { 03BC }
31648      { 039D } { 03BD }
31649      { 039E } { 03BE }
31650      { 039F } { 03BF }
31651      { 03A0 } { 03C0 }
31652      { 03A1 } { 03C1 }
31653      { 03A3 } { 03C3 }
31654      { 03A4 } { 03C4 }
31655      { 03A5 } { 03C5 }
31656      { 03A6 } { 03C6 }
31657      { 03A7 } { 03C7 }
31658      { 03A8 } { 03C8 }
31659      { 03A9 } { 03C9 }
31660      { 03AA } { 03CA }
31661      { 03AB } { 03CB }
31662      { 038C } { 03CC }
31663      { 038E } { 03CD }
31664      { 038F } { 03CE }
31665      { 03CF } { 03D7 }
31666      { 03D8 } { 03D9 }
31667      { 03DA } { 03DB }
31668      { 03DC } { 03DD }
31669      { 03DE } { 03DF }
31670      { 03E0 } { 03E1 }
31671      { 03E2 } { 03E3 }
31672      { 03E4 } { 03E5 }
31673      { 03E6 } { 03E7 }
31674      { 03E8 } { 03E9 }
31675      { 03EA } { 03EB }
31676      { 03EC } { 03ED }
31677      { 03EE } { 03EF }
31678      { 03F9 } { 03F2 }
31679      { 037F } { 03F3 }
31680      { 03F7 } { 03F8 }
31681      { 03FA } { 03FB }
31682      \q_recursion_tail ?

```

```
31683 \q_recursion_stop
```

Odds and ends for Greek; mainly symbols that are for compatibility, but also things like the terminal sigma. Almost all are uppercase mappings, but there is one that is not!

```
31684 \cs_set_protected:Npn \__text_tmp:w #1#2#3
31685 {
31686   \group_begin:
31687   \cs_set_protected:Npn \__text_tmp:w ##1##2##3##4##5##6##7##8
31688   {
31689     \tl_const:cx
31690     {
31691       c__text_ #3 case_
31692       \char_generate:nn {##1} { 12 }
31693       \char_generate:nn {##2} { 12 }
31694       _tl
31695     }
31696     {
31697       \exp_after:wN \exp_after:wN \exp_after:wN
31698       \exp_not:N \char_generate:nn {##5} { 13 }
31699       \exp_after:wN \exp_after:wN \exp_after:wN
31700       \exp_not:N \char_generate:nn {##6} { 13 }
31701     }
31702   }
31703   \use:x
31704   {
31705     \__text_tmp:w
31706     \char_to_utfviii_bytes:n { "#1 }
31707     \char_to_utfviii_bytes:n { "#2 }
31708   }
31709   \group_end:
31710 }
31711 \__text_tmp:w { 0345 } { 0399 } { upper }
31712 \__text_tmp:w { 03C2 } { 03A3 } { upper }
31713 \__text_tmp:w { 03D0 } { 0392 } { upper }
31714 \__text_tmp:w { 03D1 } { 0398 } { upper }
31715 \__text_tmp:w { 03D5 } { 03A6 } { upper }
31716 \__text_tmp:w { 03D6 } { 03A0 } { upper }
31717 \__text_tmp:w { 03F0 } { 039A } { upper }
31718 \__text_tmp:w { 03F1 } { 03A1 } { upper }
31719 \__text_tmp:w { 03F4 } { 03B8 } { lower }
31720 \__text_tmp:w { 03F5 } { 0395 } { upper }
```

Odds and ends that are not simple one-to-one mappings. These are still two-octet code points.

```
31721 \cs_set_protected:Npn \__text_tmp:w #1#2#3
31722 {
31723   \group_begin:
31724   \cs_set_protected:Npn \__text_tmp:w ##1##2##3##4
31725   {
31726     \tl_const:cn
31727     {
31728       c__text_ #3 case_
31729       \char_generate:nn {##1} { 12 }
31730       \char_generate:nn {##2} { 12 }
31731       _tl
```

```

31732         }
31733         {#2}
31734     }
31735     \use:x
31736     { \__text_tmp:w \char_to_utfviii_bytes:n { "#1 } }
31737     \group_end:
31738 }
31739 \__text_tmp:w { 00DF } { SS } { upper }
31740 \__text_tmp:w { 00DF } { Ss } { title }
31741 \__text_tmp:w { 0131 } { I } { upper }

```

Greek support: the three-octet code points.

```

31742 \cs_set_protected:Npn \__text_tmp:nnnnnn #1#2#3#4#5#6#7
31743 {
31744     \tl_const:cx
31745     {
31746         c__text_ #1 case_
31747         \char_generate:nn {#2} { 12 }
31748         \char_generate:nn {#3} { 12 }
31749         \char_generate:nn {#4} { 12 }
31750         _tl
31751     }
31752     {
31753         \exp_after:wN \exp_after:wN \exp_after:wN
31754         \exp_not:N \char_generate:nn {#5} { 13 }
31755         \exp_after:wN \exp_after:wN \exp_after:wN
31756         \exp_not:N \char_generate:nn {#6} { 13 }
31757         \exp_after:wN \exp_after:wN \exp_after:wN
31758         \exp_not:N \char_generate:nn {#7} { 13 }
31759     }
31760 }
31761 \cs_set_protected:Npn \__text_tmp:w #1#2#3#4#5#6#7#8
31762 {
31763     \tl_const:cx
31764     {
31765         c__text_lowercase_
31766         \char_generate:nn {#1} { 12 }
31767         \char_generate:nn {#2} { 12 }
31768         \char_generate:nn {#3} { 12 }
31769         _tl
31770     }
31771     {
31772         \exp_after:wN \exp_after:wN \exp_after:wN
31773         \exp_not:N \char_generate:nn {#5} { 13 }
31774         \exp_after:wN \exp_after:wN \exp_after:wN
31775         \exp_not:N \char_generate:nn {#6} { 13 }
31776         \exp_after:wN \exp_after:wN \exp_after:wN
31777         \exp_not:N \char_generate:nn {#7} { 13 }
31778     }
31779     \__text_tmp:nnnnnn { upper } {#5} {#6} {#7} {#1} {#2} {#3}
31780     \__text_tmp:nnnnnn { title } {#5} {#6} {#7} {#1} {#2} {#3}
31781 }
31782 \__text_loop:nn
31783 { 1F08 } { 1F00 }
31784 { 1F09 } { 1F01 }

```

31785	{ 1FOA }	{ 1F02 }
31786	{ 1FOB }	{ 1F03 }
31787	{ 1FOC }	{ 1F04 }
31788	{ 1FOD }	{ 1F05 }
31789	{ 1FOE }	{ 1F06 }
31790	{ 1F0F }	{ 1F07 }
31791	{ 1F18 }	{ 1F10 }
31792	{ 1F19 }	{ 1F11 }
31793	{ 1F1A }	{ 1F12 }
31794	{ 1F1B }	{ 1F13 }
31795	{ 1F1C }	{ 1F14 }
31796	{ 1F1D }	{ 1F15 }
31797	{ 1F28 }	{ 1F20 }
31798	{ 1F29 }	{ 1F21 }
31799	{ 1F2A }	{ 1F22 }
31800	{ 1F2B }	{ 1F23 }
31801	{ 1F2C }	{ 1F24 }
31802	{ 1F2D }	{ 1F25 }
31803	{ 1F2E }	{ 1F26 }
31804	{ 1F2F }	{ 1F27 }
31805	{ 1F38 }	{ 1F30 }
31806	{ 1F39 }	{ 1F31 }
31807	{ 1F3A }	{ 1F32 }
31808	{ 1F3B }	{ 1F33 }
31809	{ 1F3C }	{ 1F34 }
31810	{ 1F3D }	{ 1F35 }
31811	{ 1F3E }	{ 1F36 }
31812	{ 1F3F }	{ 1F37 }
31813	{ 1F48 }	{ 1F40 }
31814	{ 1F49 }	{ 1F41 }
31815	{ 1F4A }	{ 1F42 }
31816	{ 1F4B }	{ 1F43 }
31817	{ 1F4C }	{ 1F44 }
31818	{ 1F4D }	{ 1F45 }
31819	{ 1F59 }	{ 1F51 }
31820	{ 1F5B }	{ 1F53 }
31821	{ 1F5D }	{ 1F55 }
31822	{ 1F5F }	{ 1F57 }
31823	{ 1F68 }	{ 1F60 }
31824	{ 1F69 }	{ 1F61 }
31825	{ 1F6A }	{ 1F62 }
31826	{ 1F6B }	{ 1F63 }
31827	{ 1F6C }	{ 1F64 }
31828	{ 1F6D }	{ 1F65 }
31829	{ 1F6E }	{ 1F66 }
31830	{ 1F6F }	{ 1F67 }
31831	{ 1FBA }	{ 1F70 }
31832	{ 1FBB }	{ 1F71 }
31833	{ 1FC8 }	{ 1F72 }
31834	{ 1FC9 }	{ 1F73 }
31835	{ 1FCA }	{ 1F74 }
31836	{ 1FCB }	{ 1F75 }
31837	{ 1FDA }	{ 1F76 }
31838	{ 1FDB }	{ 1F77 }


```

31839      { 1FF8 } { 1F78 }
31840      { 1FF9 } { 1F79 }
31841      { 1FEA } { 1F7A }
31842      { 1FEB } { 1F7B }
31843      { 1FFA } { 1F7C }
31844      { 1FFB } { 1F7D }
31845      { 1F88 } { 1F80 }
31846      { 1F89 } { 1F81 }
31847      { 1F8A } { 1F82 }
31848      { 1F8B } { 1F83 }
31849      { 1F8C } { 1F84 }
31850      { 1F8D } { 1F85 }
31851      { 1F8E } { 1F86 }
31852      { 1F8F } { 1F87 }
31853      { 1F98 } { 1F90 }
31854      { 1F99 } { 1F91 }
31855      { 1F9A } { 1F92 }
31856      { 1F9B } { 1F93 }
31857      { 1F9C } { 1F94 }
31858      { 1F9D } { 1F95 }
31859      { 1F9E } { 1F96 }
31860      { 1F9F } { 1F97 }
31861      { 1FA8 } { 1FA0 }
31862      { 1FA9 } { 1FA1 }
31863      { 1FAA } { 1FA2 }
31864      { 1FAB } { 1FA3 }
31865      { 1FAC } { 1FA4 }
31866      { 1FAD } { 1FA5 }
31867      { 1FAE } { 1FA6 }
31868      { 1FAF } { 1FA7 }
31869      { 1FB8 } { 1FB0 }
31870      { 1FB9 } { 1FB1 }
31871      { 1FBC } { 1FB3 }
31872      { 1FCC } { 1FC3 }
31873      { 1FD8 } { 1FD0 }
31874      { 1FD9 } { 1FD1 }
31875      { 1FE8 } { 1FE0 }
31876      { 1FE9 } { 1FE1 }
31877      { 1FEC } { 1FE5 }
31878      { 1FFC } { 1FF3 }
31879      \q_recursion_tail ?
31880      \q_recursion_stop

```

One three-octet special case for Greek: it also moves to two-octets!

```

31881      \cs_set_protected:Npn \__text_tmp:w #1#2#3
31882      {
31883          \group_begin:
31884          \cs_set_protected:Npn \__text_tmp:w ##1##2##3##4##5##6##7##8
31885          {
31886              \tl_const:cx
31887              {
31888                  c__text_ #3 case_
31889                  \char_generate:nn {##1} { 12 }
31890                  \char_generate:nn {##2} { 12 }
31891                  \char_generate:nn {##3} { 12 }

```

```

31892         _tl
31893     }
31894     {
31895         \exp_after:wN \exp_after:wN \exp_after:wN
31896         \exp_not:N \char_generate:nn {##5} { 13 }
31897         \exp_after:wN \exp_after:wN \exp_after:wN
31898         \exp_not:N \char_generate:nn {##6} { 13 }
31899     }
31900 }
31901 \use:x
31902 {
31903     \__text_tmp:w
31904     \char_to_utfviii_bytes:n { "#1 }
31905     \char_to_utfviii_bytes:n { "#2 }
31906 }
31907 \group_end:
31908 }
31909 \__text_tmp:w { 1FBE } { 0399 } { upper }
31910 }
31911 \group_end:

```

The (fixed) look-up mappings for letter-like control sequences.

```

31912 \group_begin:
31913 \cs_set_protected:Npn \__text_change_case_setup:NN #1#2
31914 {
31915     \quark_if_recursion_tail_stop:N #1
31916     \tl_const:cn { c__text_lowercase_ \token_to_str:N #1 _tl }
31917     { #2 }
31918     \tl_const:cn { c__text_uppercase_ \token_to_str:N #2 _tl }
31919     { #1 }
31920     \__text_change_case_setup:NN
31921 }
31922 \__text_change_case_setup:NN
31923 \AA \aa
31924 \AE \ae
31925 \DH \dh
31926 \DJ \dj
31927 \IJ \ij
31928 \L \l
31929 \NG \ng
31930 \O \o
31931 \OE \oe
31932 \SS \ss
31933 \TH \th
31934 \q_recursion_tail ?
31935 \q_recursion_stop
31936 \tl_const:cn { c__text_uppercase_ \token_to_str:N \i _tl } { I }
31937 \tl_const:cn { c__text_uppercase_ \token_to_str:N \j _tl } { J }
31938 \group_end:

```

To deal with possible encoding-specific extensions to `\@uclclist`, we check at the end of the preamble. This will therefore only apply to L^AT_EX 2_ε package mode.

```

31939 \cs_if_exist:cT { \@uclclist }
31940 {
31941     \AtBeginDocument

```

```

31942 {
31943   \group_begin:
31944   \cs_set_protected:Npn \__text_change_case_setup:Nn #1#2
31945   {
31946     \quark_if_recursion_tail_stop:N #1
31947     \tl_if_single_token:nT {#2}
31948     {
31949       \cs_if_exist:cF
31950       { c__text_uppercase_ \token_to_str:N #1 _tl }
31951       {
31952         \tl_const:cn
31953         { c__text_uppercase_ \token_to_str:N #1 _tl }
31954         { #2 }
31955       }
31956       \cs_if_exist:cF
31957       { c__text_lowercase_ \token_to_str:N #2 _tl }
31958       {
31959         \tl_const:cn
31960         { c__text_lowercase_ \token_to_str:N #2 _tl }
31961         { #1 }
31962       }
31963     }
31964     \__text_change_case_setup:Nn
31965   }
31966   \exp_after:wN \__text_change_case_setup:Nn \@uclclist
31967   \q_recursion_tail ?
31968   \q_recursion_stop
31969   \group_end:
31970 }
31971 }
31972 </package>

```

49 l3text-purify implementation

```

31973 (*package)
31974 (@@=text)

```

49.1 Purifying text

_text_if_recursion_tail_stop:N Functions to query recursion quarks.

```
31975 \__kernel_quark_new_test:N \_text_if_recursion_tail_stop:N
```

(End definition for _text_if_recursion_tail_stop:N.)

\text_purify:n

As in the other parts of the module, we start off with a standard “action” loop, with expansion applied up-front.

```

\__text_purify:n
\__text_purify_store:n
\__text_purify_store:nw
\__text_purify_end:w
\__text_purify_loop:w
\__text_purify_group:n
\__text_purify_space:w
\__text_purify_N_type:N
\__text_purify_N_type_aux:N
\__text_purify_math_search:NNN
\__text_purify_math_start:NNw
\__text_purify_math_store:n
\__text_purify_math_store:nw
\__text_purify_math_end:w
\__text_purify_math_loop:NNw
\__text_purify_math_N_type:NNN
\__text_purify_math_group:NNn
\__text_purify_math_space:NNw

```

```

31976 \cs_new:Npn \text_purify:n #1
31977 {
31978   \__kernel_exp_not:w \exp_after:wN
31979   {
31980     \exp:w
31981     \exp_args:Ne \__text_purify:n
31982     { \text_expand:n {#1} }

```

```

31983     }
31984   }
31985   \cs_new:Npn \__text_purify:n #1
31986   {
31987     \group_align_safe_begin:
31988     \__text_purify_loop:w #1
31989     \q__text_recursion_tail \q__text_recursion_stop
31990     \__text_purify_result:n { }
31991   }

```

As for expansion, collect up the tokens for future use.

```

31992 \cs_new:Npn \__text_purify_store:n #1
31993 { \__text_purify_store:nw {#1} }
31994 \cs_new:Npn \__text_purify_store:nw #1#2 \__text_purify_result:n #3
31995 { #2 \__text_purify_result:n { #3 #1 } }
31996 \cs_new:Npn \__text_purify_end:w #1 \__text_purify_result:n #2
31997 {
31998   \group_align_safe_end:
31999   \exp_end:
32000   #2
32001 }

```

The main loop is a standard “tl action”. Unlike the expansion or case changing, here any groups have to be run inline. Most of the business end is as before in the N-type token processing.

```

32002 \cs_new:Npn \__text_purify_loop:w #1 \q__text_recursion_stop
32003 {
32004   \tl_if_head_is_N_type:nTF {#1}
32005   { \__text_purify_N_type:N }
32006   {
32007     \tl_if_head_is_group:nTF {#1}
32008     { \__text_purify_group:n }
32009     { \__text_purify_space:w }
32010   }
32011   #1 \q__text_recursion_stop
32012 }
32013 \cs_new:Npn \__text_purify_group:n #1 { \__text_purify_loop:w #1 }
32014 \exp_last_unbraced:NNo \cs_new:Npn \__text_purify_space:w \c_space_tl
32015 {
32016   \__text_purify_store:n { ~ }
32017   \__text_purify_loop:w
32018 }

```

The first part of handling math mode is exactly the same as in the other functions: look for a start-of-math mode token and if found start a new loop tracking the closing token.

```

32019 \cs_new:Npn \__text_purify_N_type:N #1
32020 {
32021   \__text_if_recursion_tail_stop_do:Nn #1 { \__text_purify_end:w }
32022   \__text_purify_N_type_aux:N #1
32023 }
32024 \cs_new:Npn \__text_purify_N_type_aux:N #1
32025 {
32026   \exp_after:wN \__text_purify_math_search:NNN
32027   \exp_after:wN #1 \l_text_math_delims_tl
32028   \q__text_recursion_tail ?

```

```

32029     \q__text_recursion_stop
32030   }
32031 \cs_new:Npn \__text_purify_math_search:NNN #1#2#3
32032 {
32033   \__text_if_recursion_tail_stop_do:Nn #2
32034   { \__text_purify_math_cmd:N #1 }
32035   \token_if_eq_meaning:NNTF #1 #2
32036   {
32037     \__text_use_i_delimit_by_q_recursion_stop:nw
32038     { \__text_purify_math_start:NNw #2 #3 }
32039   }
32040   { \__text_purify_math_search:NNN #1 }
32041 }
32042 \cs_new:Npn \__text_purify_math_start:NNw #1#2#3 \q__text_recursion_stop
32043 {
32044   \__text_purify_math_loop:NNw #1#2#3 \q__text_recursion_stop
32045   \__text_purify_math_result:n { }
32046 }
32047 \cs_new:Npn \__text_purify_math_store:n #1
32048 { \__text_purify_math_store:nw {#1} }
32049 \cs_new:Npn \__text_purify_math_store:nw #1#2 \__text_purify_math_result:n #3
32050 { #2 \__text_purify_math_result:n { #3 #1 } }
32051 \cs_new:Npn \__text_purify_math_end:w #1 \__text_purify_math_result:n #2
32052 {
32053   \__text_purify_store:n { $ #2 $ }
32054   \__text_purify_loop:w #1
32055 }
32056 \cs_new:Npn \__text_purify_math_stop:Nw #1 \__text_purify_math_result:n #2
32057 {
32058   \__text_purify_store:n {#1#2}
32059   \__text_purify_end:w
32060 }
32061 \cs_new:Npn \__text_purify_math_loop:NNw #1#2#3 \q__text_recursion_stop
32062 {
32063   \tl_if_head_is_N_type:nTF {#3}
32064   { \__text_purify_math_N_type:NNN }
32065   {
32066     \tl_if_head_is_group:nTF {#3}
32067     { \__text_purify_math_group:NNn }
32068     { \__text_purify_math_space:NNw }
32069   }
32070   #1#2#3 \q__text_recursion_stop
32071 }
32072 \cs_new:Npn \__text_purify_math_N_type:NNN #1#2#3
32073 {
32074   \__text_if_recursion_tail_stop_do:Nn #3
32075   { \__text_purify_math_stop:Nw #1 }
32076   \token_if_eq_meaning:NNTF #3 #2
32077   { \__text_purify_math_end:w }
32078   {
32079     \__text_purify_math_store:n {#3}
32080     \__text_purify_math_loop:NNw #1#2
32081   }
32082 }

```

```

32083 \cs_new:Npn \__text_purify_math_group:NNn #1#2#3
32084 {
32085   \__text_purify_math_store:n { {#3} }
32086   \__text_purify_math_loop:NNw #1#2
32087 }
32088 \exp_after:wN \cs_new:Npn \exp_after:wN \__text_purify_math_space:NNw
32089 \exp_after:wN # \exp_after:wN 1
32090 \exp_after:wN # \exp_after:wN 2 \c_space_tl
32091 {
32092   \__text_purify_math_store:n { ~ }
32093   \__text_purify_math_loop:NNw #1#2
32094 }

```

Then handle math mode as an argument: same outcomes, different input syntax.

```

32095 \cs_new:Npn \__text_purify_math_cmd:N #1
32096 {
32097   \exp_after:wN \__text_purify_math_cmd:NN \exp_after:wN #1
32098   \l_text_math_arg_tl \q_text_recursion_tail \q_text_recursion_stop
32099 }
32100 \cs_new:Npn \__text_purify_math_cmd:NN #1#2
32101 {
32102   \__text_if_recursion_tail_stop_do:Nn #2
32103   { \__text_purify_replace:N #1 }
32104   \cs_if_eq:NNTF #2 #1
32105   {
32106     \__text_use_i_delimit_by_q_recursion_stop:nw
32107     { \__text_purify_math_cmd:n }
32108   }
32109   { \__text_purify_math_cmd:NN #1 }
32110 }
32111 \cs_new:Npn \__text_purify_math_cmd:n #1
32112 { \__text_purify_math_end:w \__text_purify_math_result:n {#1} }

```

For N-type tokens, we first look for a string-context replacement before anything else: this can therefore cover anything. Assuming we don't find one, check to see if we can expand control sequences: if not, they have to be dropped. We also allow for L^AT_EX 2_ε `\protect`: there's an assumption that we don't have `\protect { \oops }` or similar, but that's also in the expansion code and seems like a reasonable balance.

```

32113 \cs_new:Npn \__text_purify_replace:N #1
32114 {
32115   \bool_lazy_and:nnTF
32116   { \cs_if_exist_p:c { l__text_purify_ \token_to_str:N #1 _tl } }
32117   {
32118     \bool_lazy_or_p:nn
32119     { \token_if_cs_p:N #1 }
32120     { \token_if_active_p:N #1 }
32121   }
32122   {
32123     \exp_args:Nv \__text_purify_replace:n
32124     { l__text_purify_ \token_to_str:N #1 _tl }
32125   }
32126   {
32127     \token_if_cs:NNTF #1
32128     { \__text_purify_expand:N #1 }
32129     {

```

```

32130         \exp_args:Ne \__text_purify_store:n
32131         { \__text_token_to_explicit:N #1 }
32132         \__text_purify_loop:w
32133     }
32134 }
32135 }
32136 \cs_new:Npn \__text_purify_replace:n #1 { \__text_purify_loop:w #1 }
32137 \cs_new:Npn \__text_purify_expand:N #1
32138 {
32139     \str_if_eq:nnTF {#1} { \protect }
32140     { \__text_purify_protect:N }
32141     { \__text_purify_encoding:N #1 }
32142 }
32143 \cs_new:Npn \__text_purify_protect:N #1
32144 {
32145     \__text_if_recursion_tail_stop_do:Nn #1 { \__text_purify_end:w }
32146     \__text_purify_loop:w
32147 }

```

Handle encoding commands, as detailed for expansion.

```

32148 \cs_new:Npn \__text_purify_encoding:N #1
32149 {
32150     \bool_lazy_or:nnTF
32151     { \cs_if_eq_p:NN #1 \@current@cmd }
32152     { \cs_if_eq_p:NN #1 \@changed@cmd }
32153     { \__text_purify_encoding_escape:NN }
32154     {
32155         \__text_if_expandable:NTF #1
32156         { \exp_after:wN \__text_purify_loop:w #1 }
32157         { \__text_purify_loop:w }
32158     }
32159 }
32160 \cs_new:Npn \__text_purify_encoding_escape:NN #1#2
32161 {
32162     \__text_purify_store:n {#1}
32163     \__text_purify_loop:w
32164 }

```

(End definition for `\text_purify:n` and others. This function is documented on page 266.)

`\text_declare_purify_equivalent:Nn`

`\text_declare_purify_equivalent:Nx`

```

32165 \cs_new_protected:Npn \text_declare_purify_equivalent:Nn #1#2
32166 {
32167     \tl_clear_new:c { l__text_purify_ \token_to_str:N #1 _tl }
32168     \tl_set:cn { l__text_purify_ \token_to_str:N #1 _tl } {#2}
32169 }
32170 \cs_generate_variant:Nn \text_declare_purify_equivalent:Nn { Nx }

```

(End definition for `\text_declare_purify_equivalent:Nn`. This function is documented on page 266.)

Now pre-define a range of standard commands that need dedicated definitions in purified text. First handle font-related stuff: all of this needs to be disabled.

```

32171 \tl_map_inline:nn
32172 {
32173     \fontencoding
32174     \fontfamily

```

```

32175     \fontseries
32176     \fontshape
32177   }
32178   { \text_declare_purify_equivalent:Nn #1 { \use_none:n } }
32179 \text_declare_purify_equivalent:Nn \fontsize { \use_none:nn }
32180 \text_declare_purify_equivalent:Nn \selectfont { }
32181 \text_declare_purify_equivalent:Nn \usefont { \use_none:nnnn }
32182 \tl_map_inline:nn
32183 {
32184   \emph
32185   \text
32186   \textnormal
32187   \textrm
32188   \textsf
32189   \texttt
32190   \textbf
32191   \textmd
32192   \textit
32193   \textsl
32194   \textup
32195   \textsc
32196   \textulc
32197 }
32198 { \text_declare_purify_equivalent:Nn #1 { \use:n } }
32199 \tl_map_inline:nn
32200 {
32201   \normalfont
32202   \rmfamily
32203   \sffamily
32204   \ttfamily
32205   \bfseries
32206   \mdseries
32207   \itshape
32208   \scshape
32209   \slshape
32210   \upshape
32211   \em
32212   \Huge
32213   \LARGE
32214   \Large
32215   \footnotesize
32216   \huge
32217   \large
32218   \normalsize
32219   \scriptsize
32220   \small
32221   \tiny
32222 }
32223 { \text_declare_purify_equivalent:Nn #1 { } }
32224 \exp_args:Nc \text_declare_purify_equivalent:Nn
32225 { @protected@testopt } { \use_none:nnn }

```

Environments have to be handled by pure expansion.

__text_end_env:n


```

32226 \text_declare_purify_equivalent:Nn \begin { \use:c }
32227 \text_declare_purify_equivalent:Nn \end { \_text_end_env:n }
32228 \cs_new:Npn \_text_end_env:n #1 { \cs:w end #1 \cs_end: }

```

(End definition for `_text_end_env:n`.)

Some common symbols and similar ideas.

```

32229 \text_declare_purify_equivalent:Nn \ { }
32230 \tl_map_inline:nn
32231 { \{ \} \# \$ \% \_ }
32232 { \text_declare_purify_equivalent:Nx #1 { \cs_to_str:N #1 } }

```

Cross-referencing.

```

32233 \text_declare_purify_equivalent:Nn \label { \use_none:n }

```

Spaces.

```

32234 \group_begin:
32235 \char_set_catcode_active:N \~
32236 \use:n
32237 {
32238   \group_end:
32239   \text_declare_purify_equivalent:Nx ~ { \c_space_tl }
32240 }
32241 \text_declare_purify_equivalent:Nn \nobreakspace { ~ }
32242 \text_declare_purify_equivalent:Nn \ { ~ }
32243 \text_declare_purify_equivalent:Nn \, { ~ }

```

49.2 Accent and letter-like data for purifying text

In contrast to case changing, both 8-bit and Unicode engines need information for text purification to handle accents and letter-like functions: these all need to be removed. However, the results are of course engine-dependent.

For the letter-like commands, life is relatively easy: they are all simply added as standard exceptions. The only oddity is `\SS`, which gets converted to two letters. (At some stage an alternative version can presumably be added to `babel` or similar.)

```

32244 \bool_lazy_or:nnTF
32245 { \sys_if_engine_luatex_p: }
32246 { \sys_if_engine_xetex_p: }
32247 {
32248   \cs_set_protected:Npn \_text_loop:Nn #1#2
32249   {
32250     \quark_if_recursion_tail_stop:N #1
32251     \text_declare_purify_equivalent:Nx #1
32252     {
32253       \char_generate:nn { "#2 }
32254       { \char_value_catcode:n { "#2 } }
32255     }
32256     \_text_loop:Nn
32257   }
32258 }
32259 {
32260   \cs_set_protected:Npn \_text_loop:Nn #1#2
32261   {
32262     \quark_if_recursion_tail_stop:N #1
32263     \text_declare_purify_equivalent:Nx #1

```

```

32264         {
32265             \exp_args:Ne \__text_tmp:n
32266             { \char_to_utfviii_bytes:n { "#2 } }
32267         }
32268     \__text_loop:Nn
32269 }
32270 \cs_set:Npn \__text_tmp:n #1 { \__text_tmp:nnnn #1 }
32271 \cs_set:Npn \__text_tmp:nnnn #1#2#3#4
32272 {
32273     \exp_after:wN \exp_after:wN \exp_after:wN
32274     \exp_not:N \char_generate:nn {#1} { 13 }
32275     \exp_after:wN \exp_after:wN \exp_after:wN
32276     \exp_not:N \char_generate:nn {#2} { 13 }
32277 }
32278 }
32279 \__text_loop:Nn
32280 \AA { 00C5 }
32281 \AE { 00C6 }
32282 \DH { 00D0 }
32283 \DJ { 0110 }
32284 \IJ { 0132 }
32285 \L { 0141 }
32286 \NG { 014A }
32287 \O { 00D8 }
32288 \OE { 0152 }
32289 \TH { 00DE }
32290 \aa { 00E5 }
32291 \ae { 00E6 }
32292 \dh { 00F0 }
32293 \dj { 0111 }
32294 \i { 0131 }
32295 \j { 0237 }
32296 \ij { 0132 }
32297 \l { 0142 }
32298 \ng { 014B }
32299 \o { 00F8 }
32300 \oe { 0153 }
32301 \ss { 00DF }
32302 \th { 00FE }
32303 \q_recursion_tail ?
32304 \q_recursion_stop
32305 \text_declare_purify_equivalent:Nn \SS { SS }

```

__text_purify_accent:NN Accent LICR handling is a little more complex. Accents may exist as pre-composed codepoints or as independent glyphs. The former are all saved as single token lists, whilst for the latter the combining accent needs to be re-ordered compared to the character it applies to.

```

32306 \cs_new:Npn \__text_purify_accent:NN #1#2
32307 {
32308     \cs_if_exist:cTF
32309     { c__text_purify_ \token_to_str:N #1 _ \token_to_str:N #2 _t1 }
32310     {
32311         \exp_not:v
32312         { c__text_purify_ \token_to_str:N #1 _ \token_to_str:N #2 _t1 }

```

```

32313     }
32314     {
32315         \exp_not:n {#2}
32316         \exp_not:v { c__text_purify_ \token_to_str:N #1 _tl }
32317     }
32318 }
32319 \tl_map_inline:Nn \l_text_accents_tl
32320 { \text_declare_purify_equivalent:Nn #1 { \__text_purify_accent:NN #1 } }

```

First set up the combining accents.

```

32321 \group_begin:
32322   \cs_set_protected:Npn \__text_loop:Nn #1#2
32323   {
32324     \quark_if_recursion_tail_stop:N #1
32325     \tl_const:cx { c__text_purify_ \token_to_str:N #1 _tl }
32326     { \__text_tmp:n {#2} }
32327     \__text_loop:Nn
32328   }
32329   \bool_lazy_or:nnTF
32330   { \sys_if_engine luatex_p: }
32331   { \sys_if_engine xetex_p: }
32332   {
32333     \cs_set:Npn \__text_tmp:n #1
32334     {
32335       \char_generate:nn { "#1 }
32336       { \char_value_catcode:n { "#1 } }
32337     }
32338   }
32339   {
32340     \cs_set:Npn \__text_tmp:n #1
32341     {
32342       \exp_args:Ne \__text_tmp_aux:n
32343       { \char_to_utfviii_bytes:n { "#1 } }
32344     }
32345     \cs_set:Npn \__text_tmp_aux:n #1 { \__text_tmp:nnnn #1 }
32346     \cs_set:Npn \__text_tmp:nnnn #1#2#3#4
32347     {
32348       \exp_after:wN \exp_after:wN \exp_after:wN
32349       \exp_not:N \char_generate:nn {#1} { 13 }
32350       \exp_after:wN \exp_after:wN \exp_after:wN
32351       \exp_not:N \char_generate:nn {#2} { 13 }
32352     }
32353   }
32354   \__text_loop:Nn
32355   \‘ { 0300 }
32356   \’ { 0301 }
32357   \^ { 0302 }
32358   \~ { 0303 }
32359   \= { 0304 }
32360   \u { 0306 }
32361   \. { 0307 }
32362   \" { 0308 }
32363   \r { 030A }
32364   \H { 030B }
32365   \v { 030C }

```

```

32366 \d { 0323 }
32367 \c { 0327 }
32368 \k { 0328 }
32369 \b { 0331 }
32370 \t { 0361 }
32371 \q_recursion_tail { }
32372 \q_recursion_stop

```

Now we handle the pre-composed accents: the list here is taken from `puenc.def`. All of the precomposed cases take a single letter as their second argument. We do not try to cover the case where an accent is added to a “real” dotless-i or -j, or a æ/Æ. Rather, we assume that if the UTF-8 character is used, it will have the real accent character too.

```

32373 \cs_set_protected:Npn \__text_loop:NNn #1#2#3
32374 {
32375   \quark_if_recursion_tail_stop:N #1
32376   \tl_const:cx
32377   { c__text_purify_ \token_to_str:N #1 _ \token_to_str:N #2 _tl }
32378   { \__text_tmp:n {#3} }
32379   \__text_loop:NNn
32380 }
32381 \__text_loop:NNn
32382 \‘ A { 00C0 }
32383 \’ A { 00C1 }
32384 \^ A { 00C2 }
32385 \~ A { 00C3 }
32386 \" A { 00C4 }
32387 \r A { 00C5 }
32388 \c C { 00C7 }
32389 \‘ E { 00C8 }
32390 \’ E { 00C9 }
32391 \^ E { 00CA }
32392 \" E { 00CB }
32393 \‘ I { 00CC }
32394 \’ I { 00CD }
32395 \^ I { 00CE }
32396 \" I { 00CF }
32397 \~ N { 00D1 }
32398 \‘ O { 00D2 }
32399 \’ O { 00D3 }
32400 \^ O { 00D4 }
32401 \~ O { 00D5 }
32402 \" O { 00D6 }
32403 \‘ U { 00D9 }
32404 \’ U { 00DA }
32405 \^ U { 00DB }
32406 \" U { 00DC }
32407 \’ Y { 00DD }
32408 \‘ a { 00E0 }
32409 \’ a { 00E1 }
32410 \^ a { 00E2 }
32411 \~ a { 00E3 }
32412 \" a { 00E4 }
32413 \r a { 00E5 }
32414 \c c { 00E7 }

```

32415	\‘ e	{ 00E8 }
32416	\’ e	{ 00E9 }
32417	\^ e	{ 00EA }
32418	\" e	{ 00EB }
32419	\‘ i	{ 00EC }
32420	\‘ \i	{ 00EC }
32421	\’ i	{ 00ED }
32422	\’ \i	{ 00ED }
32423	\^ i	{ 00EE }
32424	\^ \i	{ 00EE }
32425	\" i	{ 00EF }
32426	\" \i	{ 00EF }
32427	\~ n	{ 00F1 }
32428	\‘ o	{ 00F2 }
32429	\’ o	{ 00F3 }
32430	\^ o	{ 00F4 }
32431	\~ o	{ 00F5 }
32432	\" o	{ 00F6 }
32433	\‘ u	{ 00F9 }
32434	\’ u	{ 00FA }
32435	\^ u	{ 00FB }
32436	\" u	{ 00FC }
32437	\’ y	{ 00FD }
32438	\" y	{ 00FF }
32439	\= A	{ 0100 }
32440	\= a	{ 0101 }
32441	\u A	{ 0102 }
32442	\u a	{ 0103 }
32443	\k A	{ 0104 }
32444	\k a	{ 0105 }
32445	\’ C	{ 0106 }
32446	\’ c	{ 0107 }
32447	\^ C	{ 0108 }
32448	\^ c	{ 0109 }
32449	\. C	{ 010A }
32450	\. c	{ 010B }
32451	\v C	{ 010C }
32452	\v c	{ 010D }
32453	\v D	{ 010E }
32454	\v d	{ 010F }
32455	\= E	{ 0112 }
32456	\= e	{ 0113 }
32457	\u E	{ 0114 }
32458	\u e	{ 0115 }
32459	\. E	{ 0116 }
32460	\. e	{ 0117 }
32461	\k E	{ 0118 }
32462	\k e	{ 0119 }
32463	\v E	{ 011A }
32464	\v e	{ 011B }
32465	\^ G	{ 011C }
32466	\^ g	{ 011D }
32467	\u G	{ 011E }
32468	\u g	{ 011F }

32469	\. G	{ 0120 }
32470	\. g	{ 0121 }
32471	\c G	{ 0122 }
32472	\c g	{ 0123 }
32473	\^ H	{ 0124 }
32474	\^ h	{ 0125 }
32475	\~ I	{ 0128 }
32476	\~ i	{ 0129 }
32477	\~ \i	{ 0129 }
32478	\= I	{ 012A }
32479	\= i	{ 012B }
32480	\= \i	{ 012B }
32481	\u I	{ 012C }
32482	\u i	{ 012D }
32483	\u \i	{ 012D }
32484	\k I	{ 012E }
32485	\k i	{ 012F }
32486	\k \i	{ 012F }
32487	\. I	{ 0130 }
32488	\^ J	{ 0134 }
32489	\^ j	{ 0135 }
32490	\^ \j	{ 0135 }
32491	\c K	{ 0136 }
32492	\c k	{ 0137 }
32493	\' L	{ 0139 }
32494	\' l	{ 013A }
32495	\c L	{ 013B }
32496	\c l	{ 013C }
32497	\v L	{ 013D }
32498	\v l	{ 013E }
32499	\. L	{ 013F }
32500	\. l	{ 0140 }
32501	\' N	{ 0143 }
32502	\' n	{ 0144 }
32503	\c N	{ 0145 }
32504	\c n	{ 0146 }
32505	\v N	{ 0147 }
32506	\v n	{ 0148 }
32507	\= O	{ 014C }
32508	\= o	{ 014D }
32509	\u O	{ 014E }
32510	\u o	{ 014F }
32511	\H O	{ 0150 }
32512	\H o	{ 0151 }
32513	\' R	{ 0154 }
32514	\' r	{ 0155 }
32515	\c R	{ 0156 }
32516	\c r	{ 0157 }
32517	\v R	{ 0158 }
32518	\v r	{ 0159 }
32519	\' S	{ 015A }
32520	\' s	{ 015B }
32521	\^ S	{ 015C }
32522	\^ s	{ 015D }

32523	\c S	{ 015E }
32524	\c s	{ 015F }
32525	\v S	{ 0160 }
32526	\v s	{ 0161 }
32527	\c T	{ 0162 }
32528	\c t	{ 0163 }
32529	\v T	{ 0164 }
32530	\v t	{ 0165 }
32531	\~ U	{ 0168 }
32532	\~ u	{ 0169 }
32533	\= U	{ 016A }
32534	\= u	{ 016B }
32535	\u U	{ 016C }
32536	\u u	{ 016D }
32537	\r U	{ 016E }
32538	\r u	{ 016F }
32539	\H U	{ 0170 }
32540	\H u	{ 0171 }
32541	\k U	{ 0172 }
32542	\k u	{ 0173 }
32543	\^ W	{ 0174 }
32544	\^ w	{ 0175 }
32545	\^ Y	{ 0176 }
32546	\^ y	{ 0177 }
32547	\" Y	{ 0178 }
32548	\' Z	{ 0179 }
32549	\' z	{ 017A }
32550	\. Z	{ 017B }
32551	\. z	{ 017C }
32552	\v Z	{ 017D }
32553	\v z	{ 017E }
32554	\v A	{ 01CD }
32555	\v a	{ 01CE }
32556	\v I	{ 01CF }
32557	\v \i	{ 01D0 }
32558	\v i	{ 01D0 }
32559	\v O	{ 01D1 }
32560	\v o	{ 01D2 }
32561	\v U	{ 01D3 }
32562	\v u	{ 01D4 }
32563	\v G	{ 01E6 }
32564	\v g	{ 01E7 }
32565	\v K	{ 01E8 }
32566	\v k	{ 01E9 }
32567	\k O	{ 01EA }
32568	\k o	{ 01EB }
32569	\v \j	{ 01F0 }
32570	\v j	{ 01F0 }
32571	\' G	{ 01F4 }
32572	\' g	{ 01F5 }
32573	\' N	{ 01F8 }
32574	\' n	{ 01F9 }
32575	\' \AE	{ 01FC }
32576	\' \ae	{ 01FD }

```

32577      \' \0 { 01FE }
32578      \' \o { 01FF }
32579      \v H { 021E }
32580      \v h { 021F }
32581      \. A { 0226 }
32582      \. a { 0227 }
32583      \c E { 0228 }
32584      \c e { 0229 }
32585      \. 0 { 022E }
32586      \. o { 022F }
32587      \= Y { 0232 }
32588      \= y { 0233 }
32589      \q_recursion_tail ? { }
32590      \q_recursion_stop
32591      \group_end:

```

(End definition for _text_purify_accent:NN.)

```

32592 \end{package}

```

50 l3legacy Implementation

```

32593 \begin{package}
32594 \@@=legacy
\legacy_if_p:n A friendly wrapper.
\legacy_if:nTF
32595 \prg_new_conditional:Npnn \legacy_if:n #1 { p , T , F , TF }
32596 {
32597     \exp_args:Nc \if_meaning:w { if#1 } \iftrue
32598     \prg_return_true:
32599     \else:
32600     \prg_return_false:
32601     \fi:
32602 }

```

(End definition for \legacy_if:nTF. This function is documented on page 267.)

```

32603 \end{package}

```

51 l3candidates Implementation

```

32604 \begin{package}

```

51.1 Additions to l3box

```

32605 \@@=box

```

51.1.1 Viewing part of a box

```

\box_clip:N A wrapper around the driver-dependent code.
\box_clip:c
\box_gclip:N
\box_gclip:c
32606 \cs_new_protected:Npn \box_clip:N #1
32607 { \hbox_set:Nn #1 { \_box_backend_clip:N #1 } }
32608 \cs_generate_variant:Nn \box_clip:N { c }
32609 \cs_new_protected:Npn \box_gclip:N #1
32610 { \hbox_gset:Nn #1 { \_box_backend_clip:N #1 } }
32611 \cs_generate_variant:Nn \box_gclip:N { c }

```


(End definition for `\box_clip:N` and `\box_gclip:N`. These functions are documented on page 268.)

```

\box_set_trim:Nnnnn Trimming from the left- and right-hand edges of the box is easy: kern the appropriate
\box_set_trim:cnnnn parts off each side.
\box_gset_trim:Nnnnn
\box_gset_trim:cnnnn
\__box_set_trim:NnnnnN
32612 \cs_new_protected:Npn \box_set_trim:Nnnnn #1#2#3#4#5
32613 { \__box_set_trim:NnnnnN #1 {#2} {#3} {#4} {#5} \box_set_eq:NN }
32614 \cs_generate_variant:Nn \box_set_trim:Nnnnn { c }
32615 \cs_new_protected:Npn \box_gset_trim:Nnnnn #1#2#3#4#5
32616 { \__box_set_trim:NnnnnN #1 {#2} {#3} {#4} {#5} \box_gset_eq:NN }
32617 \cs_generate_variant:Nn \box_gset_trim:Nnnnn { c }
32618 \cs_new_protected:Npn \__box_set_trim:NnnnnN #1#2#3#4#5#6
32619 {
32620   \hbox_set:Nn \l__box_internal_box
32621   {
32622     \tex_kern:D - \__box_dim_eval:n {#2}
32623     \box_use:N #1
32624     \tex_kern:D - \__box_dim_eval:n {#4}
32625   }

```

For the height and depth, there is a need to watch the baseline is respected. Material always has to stay on the correct side, so trimming has to check that there is enough material to trim. First, the bottom edge. If there is enough depth, simply set the depth, or if not move down so the result is zero depth. `\box_move_down:nn` is used in both cases so the resulting box always contains a `\lower` primitive. The internal box is used here as it allows safe use of `\box_set_dp:Nn`.

```

32626 \dim_compare:nNnTF { \box_dp:N #1 } > {#3}
32627 {
32628   \hbox_set:Nn \l__box_internal_box
32629   {
32630     \box_move_down:nn \c_zero_dim
32631     { \box_use_drop:N \l__box_internal_box }
32632   }
32633   \box_set_dp:Nn \l__box_internal_box { \box_dp:N #1 - (#3) }
32634 }
32635 {
32636   \hbox_set:Nn \l__box_internal_box
32637   {
32638     \box_move_down:nn { (#3) - \box_dp:N #1 }
32639     { \box_use_drop:N \l__box_internal_box }
32640   }
32641   \box_set_dp:Nn \l__box_internal_box \c_zero_dim
32642 }

```

Same thing, this time from the top of the box.

```

32643 \dim_compare:nNnTF { \box_ht:N \l__box_internal_box } > {#5}
32644 {
32645   \hbox_set:Nn \l__box_internal_box
32646   {
32647     \box_move_up:nn \c_zero_dim
32648     { \box_use_drop:N \l__box_internal_box }
32649   }
32650   \box_set_ht:Nn \l__box_internal_box
32651   { \box_ht:N \l__box_internal_box - (#5) }
32652 }

```

```

32653     {
32654         \hbox_set:Nn \l__box_internal_box
32655         {
32656             \box_move_up:nn { (#5) - \box_ht:N \l__box_internal_box }
32657             { \box_use_drop:N \l__box_internal_box }
32658         }
32659         \box_set_ht:Nn \l__box_internal_box \c_zero_dim
32660     }
32661     #6 #1 \l__box_internal_box
32662 }

```

(End definition for `\box_set_trim:Nnnnn`, `\box_gset_trim:Nnnnn`, and `__box_set_trim:NnnnnN`. These functions are documented on page 269.)

`\box_set_viewport:Nnnnn`
`\box_set_viewport:cnnnn`
`\box_gset_viewport:Nnnnn`
`\box_gset_viewport:cnnnn`
`__box_viewport:NnnnnN`

The same general logic as for the trim operation, but with absolute dimensions. As a result, there are some things to watch out for in the vertical direction.

```

32663 \cs_new_protected:Npn \box_set_viewport:Nnnnn #1#2#3#4#5
32664 { \__box_set_viewport:NnnnnN #1 {#2} {#3} {#4} {#5} \box_set_eq:NN }
32665 \cs_generate_variant:Nn \box_set_viewport:Nnnnn { c }
32666 \cs_new_protected:Npn \box_gset_viewport:Nnnnn #1#2#3#4#5
32667 { \__box_set_viewport:NnnnnN #1 {#2} {#3} {#4} {#5} \box_gset_eq:NN }
32668 \cs_generate_variant:Nn \box_gset_viewport:Nnnnn { c }
32669 \cs_new_protected:Npn \__box_set_viewport:NnnnnN #1#2#3#4#5#6
32670 {
32671     \hbox_set:Nn \l__box_internal_box
32672     {
32673         \tex_kern:D - \__box_dim_eval:n {#2}
32674         \box_use:N #1
32675         \tex_kern:D \__box_dim_eval:n { #4 - \box_wd:N #1 }
32676     }
32677     \dim_compare:nNnTF {#3} < \c_zero_dim
32678     {
32679         \hbox_set:Nn \l__box_internal_box
32680         {
32681             \box_move_down:nn \c_zero_dim
32682             { \box_use_drop:N \l__box_internal_box }
32683         }
32684         \box_set_dp:Nn \l__box_internal_box { - \__box_dim_eval:n {#3} }
32685     }
32686     {
32687         \hbox_set:Nn \l__box_internal_box
32688         { \box_move_down:nn {#3} { \box_use_drop:N \l__box_internal_box } }
32689         \box_set_dp:Nn \l__box_internal_box \c_zero_dim
32690     }
32691     \dim_compare:nNnTF {#5} > \c_zero_dim
32692     {
32693         \hbox_set:Nn \l__box_internal_box
32694         {
32695             \box_move_up:nn \c_zero_dim
32696             { \box_use_drop:N \l__box_internal_box }
32697         }
32698         \box_set_ht:Nn \l__box_internal_box
32699         {
32700             (#5)

```

```

32701         \dim_compare:nNnT {#3} > \c_zero_dim
32702         { - (#3) }
32703     }
32704 }
32705 {
32706     \hbox_set:Nn \l__box_internal_box
32707     {
32708         \box_move_up:nn { - \__box_dim_eval:n {#5} }
32709         { \box_use_drop:N \l__box_internal_box }
32710     }
32711     \box_set_ht:Nn \l__box_internal_box \c_zero_dim
32712 }
32713 #6 #1 \l__box_internal_box
32714 }

```

(End definition for `\box_set_viewport:Nnnnn`, `\box_gset_viewport:Nnnnn`, and `__box_viewport:NnnnnN`. These functions are documented on page 269.)

51.2 Additions to l3flag

```

32715 <@@=flag>

```

`\flag_raise_if_clear:n` It might be faster to just call the “trap” function in all cases but conceptually the function name suggests we should only run it if the flag is zero in case the “trap” made customizable in the future.

```

32716 \cs_new:Npn \flag_raise_if_clear:n #1
32717 {
32718     \if_cs_exist:w flag~#1~0 \cs_end:
32719     \else:
32720         \cs:w flag~#1 \cs_end: 0 ;
32721     \fi:
32722 }

```

(End definition for `\flag_raise_if_clear:n`. This function is documented on page 270.)

51.3 Additions to l3msg

```

32723 <@@=msg>

```

`\msg_show_eval:Nn` A short-hand used for `\int_show:n` and similar functions that passes to `\tl_show:n` the result of applying #1 (a function such as `\int_eval:n`) to the expression #2. The use of f-expansion ensures that #1 is expanded in the scope in which the show command is called, rather than in the group created by `\iow_wrap:nnnN`. This is only important for expressions involving the `\currentgrouplevel` or `\currentgrouptype`. On the other hand we want the expression to be converted to a string with the usual escape character, hence within the wrapping code.

```

32724 \cs_new_protected:Npn \msg_show_eval:Nn #1#2
32725 { \exp_args:Nf \__msg_show_eval:nnN { #1 {#2} } {#2} \tl_show:n }
32726 \cs_new_protected:Npn \msg_log_eval:Nn #1#2
32727 { \exp_args:Nf \__msg_show_eval:nnN { #1 {#2} } {#2} \tl_log:n }
32728 \cs_new_protected:Npn \__msg_show_eval:nnN #1#2#3 { #3 { #2 = #1 } }

```

(End definition for `\msg_show_eval:Nn`, `\msg_log_eval:Nn`, and `__msg_show_eval:nnN`. These functions are documented on page 271.)

```

\msg_show_item:n
\msg_show_item_unbraced:n
\msg_show_item:nn
\msg_show_item_unbraced:nn

```

Each item in the variable is formatted using one of the following functions. We cannot use `\\` and so on because these short-hands cannot be used inside the arguments of messages, only when defining the messages.

```

32729 \cs_new:Npx \msg_show_item:n #1
32730 { \iow_newline: > ~ \c_space_tl \exp_not:N \tl_to_str:n { {#1} } }
32731 \cs_new:Npx \msg_show_item_unbraced:n #1
32732 { \iow_newline: > ~ \c_space_tl \exp_not:N \tl_to_str:n {#1} }
32733 \cs_new:Npx \msg_show_item:nn #1#2
32734 {
32735   \iow_newline: > \use:nn { ~ } { ~ }
32736   \exp_not:N \tl_to_str:n { {#1} }
32737   \use:nn { ~ } { ~ } => \use:nn { ~ } { ~ }
32738   \exp_not:N \tl_to_str:n { {#2} }
32739 }
32740 \cs_new:Npx \msg_show_item_unbraced:nn #1#2
32741 {
32742   \iow_newline: > \use:nn { ~ } { ~ }
32743   \exp_not:N \tl_to_str:n {#1}
32744   \use:nn { ~ } { ~ } => \use:nn { ~ } { ~ }
32745   \exp_not:N \tl_to_str:n {#2}
32746 }

```

(End definition for `\msg_show_item:n` and others. These functions are documented on page 271.)

51.4 Additions to `l3prg`

```

32747 <@@=bool>

```

```

\bool_set_inverse:N
\bool_set_inverse:c
\bool_gset_inverse:N
\bool_gset_inverse:c

```

Set to false or true locally or globally.

```

32748 \cs_new_protected:Npn \bool_set_inverse:N #1
32749 { \bool_if:NTF #1 { \bool_set_false:N } { \bool_set_true:N } #1 }
32750 \cs_generate_variant:Nn \bool_set_inverse:N { c }
32751 \cs_new_protected:Npn \bool_gset_inverse:N #1
32752 { \bool_if:NTF #1 { \bool_gset_false:N } { \bool_gset_true:N } #1 }
32753 \cs_generate_variant:Nn \bool_gset_inverse:N { c }

```

(End definition for `\bool_set_inverse:N` and `\bool_gset_inverse:N`. These functions are documented on page 271.)

```

\s__bool_mark
\s__bool_stop

```

Internal scan marks.

```

32754 \scan_new:N \s__bool_mark
32755 \scan_new:N \s__bool_stop

```

(End definition for `\s__bool_mark` and `\s__bool_stop`.)

```

\bool_case_true:n
\bool_case_true:nTF
\bool_case_false:n
\bool_case_false:nTF
  \__bool_case:NnTF
  \__bool_case_true:w
  \__bool_case_false:w
  \__bool_case_end:nw

```

For boolean cases the overall idea is the same as for `\tl_case:nn(TF)` as described in `l3tl`.

```

32756 \cs_new:Npn \bool_case_true:nTF
32757 { \exp:w \__bool_case:NnTF \c_true_bool }
32758 \cs_new:Npn \bool_case_true:nT #1#2
32759 { \exp:w \__bool_case:NnTF \c_true_bool {#1} {#2} { } }
32760 \cs_new:Npn \bool_case_true:nF #1
32761 { \exp:w \__bool_case:NnTF \c_true_bool {#1} { } }
32762 \cs_new:Npn \bool_case_true:n #1
32763 { \exp:w \__bool_case:NnTF \c_true_bool {#1} { } { } }

```

```

32764 \cs_new:Npn \bool_case_false:nTF
32765   { \exp:w \__bool_case:NnTF \c_false_bool }
32766 \cs_new:Npn \bool_case_false:nT #1#2
32767   { \exp:w \__bool_case:NnTF \c_false_bool {#1} {#2} { } }
32768 \cs_new:Npn \bool_case_false:nF #1
32769   { \exp:w \__bool_case:NnTF \c_false_bool {#1} { } }
32770 \cs_new:Npn \bool_case_false:n #1
32771   { \exp:w \__bool_case:NnTF \c_false_bool {#1} { } { } }
32772 \cs_new:Npn \__bool_case:NnTF #1#2#3#4
32773   {
32774     \bool_if:NTF #1 \__bool_case_true:w \__bool_case_false:w
32775     #2 #1 { } \s__bool_mark {#3} \s__bool_mark {#4} \s__bool_stop
32776   }
32777 \cs_new:Npn \__bool_case_true:w #1#2
32778   {
32779     \bool_if:nTF {#1}
32780     { \__bool_case_end:nw {#2} }
32781     { \__bool_case_true:w }
32782   }
32783 \cs_new:Npn \__bool_case_false:w #1#2
32784   {
32785     \bool_if:nTF {#1}
32786     { \__bool_case_false:w }
32787     { \__bool_case_end:nw {#2} }
32788   }
32789 \cs_new:Npn \__bool_case_end:nw #1#2#3 \s__bool_mark #4#5 \s__bool_stop
32790   { \exp_end: #1 #4 }

```

(End definition for `\bool_case_true:nTF` and others. These functions are documented on page 272.)

51.5 Additions to `l3prop`

```

32791 <@@=prop>

```

`__prop_use_i_delimit_by_s_stop:nw` Functions to gobble up to a scan mark.

```

32792 \cs_new:Npn \__prop_use_i_delimit_by_s_stop:nw #1 #2 \s__prop_stop {#1}

```

(End definition for `__prop_use_i_delimit_by_s_stop:nw`.)

`\prop_rand_key_value:N` Contrarily to `clist`, `seq` and `tl`, there is no function to get an item of a `prop` given an integer between 1 and the number of items, so we write the appropriate code. There is no bounds checking because `\int_rand:nn` is always within bounds. The initial `\int_value:w` is stopped by the first `\s__prop` in #1.

`\prop_rand_key_value:c`
`__prop_rand_item:w`

```

32793 \cs_new:Npn \prop_rand_key_value:N #1
32794   {
32795     \prop_if_empty:NF #1
32796     {
32797       \exp_after:wN \__prop_rand_item:w
32798       \int_value:w \int_rand:nn { 1 } { \prop_count:N #1 }
32799       #1 \s__prop_stop
32800     }
32801   }
32802 \cs_generate_variant:Nn \prop_rand_key_value:N { c }
32803 \cs_new:Npn \__prop_rand_item:w #1 \s__prop \__prop_pair:wn #2 \s__prop #3
32804   {

```

```

32805 \int_compare:nNnF {#1} > 1
32806 { \__prop_use_i_delimit_by_s_stop:nw { \exp_not:n { {#2} {#3} } } }
32807 \exp_after:wN \__prop_rand_item:w
32808 \int_value:w \int_eval:n { #1 - 1 } \s__prop
32809 }

```

(End definition for `\prop_rand_key_value:N` and `__prop_rand_item:w`. This function is documented on page 272.)

51.6 Additions to l3seq

```

32810 <@@=seq>

```

```

\seq_mapthread_function:NNN
\seq_mapthread_function:NcN
\seq_mapthread_function:cNN
\seq_mapthread_function:ccN
  \__seq_mapthread_function:wNN
  \__seq_mapthread_function:wNw
  \__seq_mapthread_function:Nnnwnn

```

The idea is to first expand both sequences, adding the usual `{ ? \prg_break: } { }` to the end of each one. This is most conveniently done in two steps using an auxiliary function. The mapping then throws away the first tokens of #2 and #5, which for items in both sequences are `\s__seq __seq_item:n`. The function to be mapped are then be applied to the two entries. When the code hits the end of one of the sequences, the break material stops the entire loop and tidy up. This avoids needing to find the count of the two sequences, or worrying about which is longer.

```

32811 \cs_new:Npn \seq_mapthread_function:NNN #1#2#3
32812 { \exp_after:wN \__seq_mapthread_function:wNN #2 \s__seq_stop #1 #3 }
32813 \cs_new:Npn \__seq_mapthread_function:wNN \s__seq #1 \s__seq_stop #2#3
32814 {
32815   \exp_after:wN \__seq_mapthread_function:wNw #2 \s__seq_stop #3
32816   #1 { ? \prg_break: } { }
32817   \prg_break_point:
32818 }
32819 \cs_new:Npn \__seq_mapthread_function:wNw \s__seq #1 \s__seq_stop #2
32820 {
32821   \__seq_mapthread_function:Nnnwnn #2
32822   #1 { ? \prg_break: } { }
32823   \s__seq_stop
32824 }
32825 \cs_new:Npn \__seq_mapthread_function:Nnnwnn #1#2#3#4 \s__seq_stop #5#6
32826 {
32827   \use_none:n #2
32828   \use_none:n #5
32829   #1 {#3} {#6}
32830   \__seq_mapthread_function:Nnnwnn #1 #4 \s__seq_stop
32831 }
32832 \cs_generate_variant:Nn \seq_mapthread_function:NNN { Nc , c , cc }

```

(End definition for `\seq_mapthread_function:NNN` and others. This function is documented on page 273.)

```

\seq_set_filter:NNN
\seq_gset_filter:NNN
\__seq_set_filter:NNNn

```

Similar to `\seq_map_inline:Nn`, without a `\prg_break_point:` because the user's code is performed within the evaluation of a boolean expression, and skipping out of that would break horribly. The `__seq_wrap_item:n` function inserts the relevant `__seq_item:n` without expansion in the input stream, hence in the x-expanding assignment.

```

32833 \cs_new_protected:Npn \seq_set_filter:NNN
32834 { \__seq_set_filter:NNNn \__kernel_tl_set:Nx }
32835 \cs_new_protected:Npn \seq_gset_filter:NNN
32836 { \__seq_set_filter:NNNn \__kernel_tl_gset:Nx }

```

```

32837 \cs_new_protected:Npn \__seq_set_filter:NNNn #1#2#3#4
32838 {
32839   \__seq_push_item_def:n { \bool_if:nT {#4} { \__seq_wrap_item:n {##1} } }
32840   #1 #2 { #3 }
32841   \__seq_pop_item_def:
32842 }

```

(End definition for `\seq_set_filter:NNn`, `\seq_gset_filter:NNn`, and `__seq_set_filter:NNNn`. These functions are documented on page 273.)

`\seq_set_from_inline_x:Nnn`
`\seq_gset_from_inline_x:Nnn`
`__seq_set_from_inline_x:NNnn`

Set `__seq_item:n` then map it using the loop code.

```

32843 \cs_new_protected:Npn \seq_set_from_inline_x:Nnn
32844 { \__seq_set_from_inline_x:NNnn \__kernel_tl_set:Nx }
32845 \cs_new_protected:Npn \seq_gset_from_inline_x:Nnn
32846 { \__seq_set_from_inline_x:NNnn \__kernel_tl_gset:Nx }
32847 \cs_new_protected:Npn \__seq_set_from_inline_x:NNnn #1#2#3#4
32848 {
32849   \__seq_push_item_def:n { \exp_not:N \__seq_item:n {#4} }
32850   #1 #2 { \s__seq #3 \__seq_item:n }
32851   \__seq_pop_item_def:
32852 }

```

(End definition for `\seq_set_from_inline_x:Nnn`, `\seq_gset_from_inline_x:Nnn`, and `__seq_set_from_inline_x:NNnn`. These functions are documented on page 273.)

`\seq_set_from_function:NnN`
`\seq_gset_from_function:NnN`

Reuse `\seq_set_from_inline_x:Nnn`.

```

32853 \cs_new_protected:Npn \seq_set_from_function:NnN #1#2#3
32854 { \seq_set_from_inline_x:Nnn #1 {#2} { #3 {##1} } }
32855 \cs_new_protected:Npn \seq_gset_from_function:NnN #1#2#3
32856 { \seq_gset_from_inline_x:Nnn #1 {#2} { #3 {##1} } }

```

(End definition for `\seq_set_from_function:NnN` and `\seq_gset_from_function:NnN`. These functions are documented on page 273.)

51.7 Additions to l3sys

```

32857 <@@=sys>

```

`\c_sys_engine_version_str`

Various different engines, various different ways to extract the data!

```

32858 \str_const:Nx \c_sys_engine_version_str
32859 {
32860   \str_case:on \c_sys_engine_str
32861   {
32862     { pdftex }
32863     {
32864       \fp_eval:n { round(\int_use:N \tex_pdftexversion:D / 100 , 2) }
32865       .
32866       \tex_pdftexrevision:D
32867     }
32868   { ptex }
32869   {
32870     \cs_if_exist:NT \tex_ptexversion:D
32871     {
32872       p
32873       \int_use:N \tex_ptexversion:D

```

```

32874         .
32875         \int_use:N \tex_ptexminorversion:D
32876         \tex_ptexrevision:D
32877         -
32878         \int_use:N \tex_epTeXversion:D
32879     }
32880 }
32881 { luatex }
32882 {
32883     \fp_eval:n { round(\int_use:N \tex_luatexversion:D / 100, 2) }
32884     .
32885     \tex_luatexrevision:D
32886 }
32887 { uptex }
32888 {
32889     \cs_if_exist:NT \tex_ptexversion:D
32890     {
32891         p
32892         \int_use:N \tex_ptexversion:D
32893         .
32894         \int_use:N \tex_ptexminorversion:D
32895         \tex_ptexrevision:D
32896         -
32897         u
32898         \int_use:N \tex_uptexversion:D
32899         \tex_uptexrevision:D
32900         -
32901         \int_use:N \tex_epTeXversion:D
32902     }
32903 }
32904 { xetex }
32905 {
32906     \int_use:N \tex_XeTeXversion:D
32907     \tex_XeTeXrevision:D
32908 }
32909 }
32910 }

```

(End definition for `\c_sys_engine_version_str`. This variable is documented on page 274.)

51.8 Additions to l3file

```

32911 (@@=ior)

```

\ior_shell_open:Nn Actually much easier than either the standard open or input versions! When calling `__kernel_ior_open:Nn` the file the pipe is added to signal a shell command, but the quotes are not added yet—they are added later by `__kernel_file_name_quote:n`.

```

\__ior_shell_open:nN
32912 \cs_new_protected:Npn \ior_shell_open:Nn #1#2
32913 {
32914     \sys_if_shell:TF
32915     { \exp_args:No \__ior_shell_open:nN { \tl_to_str:n {#2} } #1 }
32916     { \__kernel_msg_error:nn { kernel } { pipe-failed } }
32917 }
32918 \cs_new_protected:Npn \__ior_shell_open:nN #1#2

```



```

32919 {
32920   \tl_if_in:nnTF {#1} { " }
32921   {
32922     \__kernel_msg_error:nnx
32923     { kernel } { quote-in-shell } {#1}
32924   }
32925   { \__kernel_ior_open:Nn #2 { |#1 } }
32926 }
32927 \__kernel_msg_new:nnnn { kernel } { pipe-failed }
32928 { Cannot~run~piped~system~commands. }
32929 {
32930   LaTeX~tried~to~call~a~system~process~but~this~was~not~possible.\\
32931   Try~the~"--shell-escape"~(or~"--enable-pipes")~option.
32932 }

```

(End definition for `\ior_shell_open:Nn` and `__ior_shell_open:nN`. This function is documented on page 270.)

51.9 Additions to l3tl

51.9.1 Building a token list

```

32933 <@@=tl>

```

Between `\tl_build_begin:N <tl var>` and `\tl_build_end:N <tl var>`, the `<tl var>` has the structure

```

\exp_end: ... \exp_end: \__tl_build_last:NNn <assignment> <next tl>
{<left>} <right>

```

where `<right>` is not braced. The “data” it represents is `<left>` followed by the “data” of `<next tl>` followed by `<right>`. The `<next tl>` is a token list variable whose name is that of `<tl var>` followed by `'`. There are between 0 and 4 `\exp_end:` to keep track of when `<left>` and `<right>` should be put into the `<next tl>`. The `<assignment>` is `\cs_set_nopar:Npx` if the variable is local, and `\cs_gset_nopar:Npx` if it is global.

`\tl_build_begin:N` First construct the `<next tl>`: using a prime here conflicts with the usual `expl3` convention but we need a name that can be derived from `#1` without any external data such as a counter. Empty that `<next tl>` and setup the structure. The local and global versions `\tl_build_gbegin:N` only differ by a single function `\cs_(g)set_nopar:Npx` used for all assignments: this is important because only that function is stored in the `<tl var>` and `<next tl>` for subsequent assignments. In principle `__tl_build_begin:NNN` could use `\tl_(g)clear_new:N` to empty `#1` and make sure it is defined, but logging the definition does not seem useful so we just do `#3 #1 {}` to clear it locally or globally as appropriate.

```

\__tl_build_begin:NN
\__tl_build_begin:NNN

```

```

32934 \cs_new_protected:Npn \tl_build_begin:N #1
32935 { \__tl_build_begin:NN \cs_set_nopar:Npx #1 }
32936 \cs_new_protected:Npn \tl_build_gbegin:N #1
32937 { \__tl_build_begin:NN \cs_gset_nopar:Npx #1 }
32938 \cs_new_protected:Npn \__tl_build_begin:NN #1#2
32939 { \exp_args:Nc \__tl_build_begin:NNN { \cs_to_str:N #2 ' } #2 #1 }
32940 \cs_new_protected:Npn \__tl_build_begin:NNN #1#2#3
32941 {
32942   #3 #1 { }
32943   #3 #2
32944   {

```

```

32945         \exp_not:n { \exp_end: \exp_end: \exp_end: \exp_end: }
32946         \exp_not:n { \__tl_build_last:NNn #3 #1 { } }
32947     }
32948 }

```

(End definition for `\tl_build_begin:N` and others. These functions are documented on page 275.)

`\tl_build_clear:N` The `begin` and `gbegin` functions already clear enough to make the token list variable effectively empty. Eventually the `begin` and `gbegin` functions should check that `#1` is empty or undefined, while the `clear` and `gclear` functions ought to empty `#1`, `#1'` and so on, similar to `\tl_build_end:N`. This only affects memory usage.

```

32949 \cs_new_eq:NN \tl_build_clear:N \tl_build_begin:N
32950 \cs_new_eq:NN \tl_build_gclear:N \tl_build_gbegin:N

```

(End definition for `\tl_build_clear:N` and `\tl_build_gclear:N`. These functions are documented on page 275.)

`\tl_build_put_right:Nn` Similar to `\tl_put_right:Nn`, but apply `\exp:w` to `#1`. Most of the time this just removes one `\exp_end:`. When there are none left, `__tl_build_last:NNn` is expanded instead.

`\tl_build_put_right:Nx` It resets the definition of the `\tl var` by ending the `\exp_not:n` and the definition early.

`\tl_build_gput_right:Nn` Then it makes sure the `\next tl` (its argument `#1`) is set-up and starts a new definition.

`\tl_build_gput_right:Nx` Then `__tl_build_put:nn` and `__tl_build_put:nw` place the `\left` part of the original `\tl var` as appropriate for the definition of the `\next tl` (the `\right` part is left in the right place without ever becoming a macro argument). We use `\exp_after:wN` rather than some `\exp_args:No` to avoid reading arguments that are likely very long token lists. We use `\cs_(g)set_nopar:Npx` rather than `\tl_(g)set:Nx` partly for the same reason and partly because the assignments are interrupted by brace tricks, which implies that the assignment does not simply set the token list to an x-expansion of the second argument.

```

32951 \cs_new_protected:Npn \tl_build_put_right:Nn #1#2
32952 {
32953     \cs_set_nopar:Npx #1
32954     { \exp_after:wN \exp_not:n \exp_after:wN { \exp:w #1 #2 } }
32955 }
32956 \cs_new_protected:Npn \tl_build_put_right:Nx #1#2
32957 {
32958     \cs_set_nopar:Npx #1
32959     { \exp_after:wN \exp_not:n \exp_after:wN { \exp:w #1 } #2 }
32960 }
32961 \cs_new_protected:Npn \tl_build_gput_right:Nn #1#2
32962 {
32963     \cs_gset_nopar:Npx #1
32964     { \exp_after:wN \exp_not:n \exp_after:wN { \exp:w #1 #2 } }
32965 }
32966 \cs_new_protected:Npn \tl_build_gput_right:Nx #1#2
32967 {
32968     \cs_gset_nopar:Npx #1
32969     { \exp_after:wN \exp_not:n \exp_after:wN { \exp:w #1 } #2 }
32970 }
32971 \cs_new_protected:Npn \__tl_build_last:NNn #1#2
32972 {
32973     \if_false: { { \fi:
32974         \exp_end: \exp_end: \exp_end: \exp_end: \exp_end:
32975         \__tl_build_last:NNn #1 #2 { }

```

```

32976     }
32977   }
32978   \if_meaning:w \c_empty_tl #2
32979   \__tl_build_begin:NN #1 #2
32980   \fi:
32981   #1 #2
32982   {
32983     \exp_after:wN \exp_not:n \exp_after:wN
32984     {
32985       \exp:w \if_false: } } \fi:
32986       \exp_after:wN \__tl_build_put:nn \exp_after:wN {#2}
32987   }
32988   \cs_new_protected:Npn \__tl_build_put:nn #1#2 { \__tl_build_put:nw {#2} #1 }
32989   \cs_new_protected:Npn \__tl_build_put:nw #1#2 \__tl_build_last:NNn #3#4#5
32990   { #2 \__tl_build_last:NNn #3 #4 { #1 #5 } }

```

(End definition for `\tl_build_put_right:Nn` and others. These functions are documented on page 276.)

`\tl_build_put_left:Nn` See `\tl_build_put_right:Nn` for all the machinery. We could easily provide `\tl_build_put_left_right:NNn`, by just add the `\right` material after the `{\left}` in the x-expanding assignment.

`\tl_build_put_left:Nx`

`\tl_build_gput_left:Nn`

`\tl_build_gput_left:Nx`

`__tl_build_put_left:NNn`

```

32991 \cs_new_protected:Npn \tl_build_put_left:Nn #1
32992 { \__tl_build_put_left:NNn \cs_set_nopar:Npx #1 }
32993 \cs_generate_variant:Nn \tl_build_put_left:Nn { Nx }
32994 \cs_new_protected:Npn \tl_build_gput_left:Nn #1
32995 { \__tl_build_put_left:NNn \cs_gset_nopar:Npx #1 }
32996 \cs_generate_variant:Nn \tl_build_gput_left:Nn { Nx }
32997 \cs_new_protected:Npn \__tl_build_put_left:NNn #1#2#3
32998 {
32999   #1 #2
33000   {
33001     \exp_after:wN \exp_not:n \exp_after:wN
33002     {
33003       \exp:w \exp_after:wN \__tl_build_put:nn
33004       \exp_after:wN {#2} {#3}
33005     }
33006   }
33007 }

```

(End definition for `\tl_build_put_left:Nn`, `\tl_build_gput_left:Nn`, and `__tl_build_put_left:NNn`. These functions are documented on page 276.)

`\tl_build_get:NN` The idea is to expand the `\tl var` then the `\next tl` and so on, all within an x-expanding assignment, and wrap as appropriate in `\exp_not:n`. The various `\left` parts are left in the assignment as we go, which enables us to expand the `\next tl` at the right place. The various `\right` parts are eventually picked up in one last `\exp_not:n`, with a brace trick to wrap all the `\right` parts together.

`__tl_build_get:NNN`

`__tl_build_get:w`

`__tl_build_get_end:w`

```

33008 \cs_new_protected:Npn \tl_build_get:NN
33009 { \__tl_build_get:NNN \__kernel_tl_set:Nx }
33010 \cs_new_protected:Npn \__tl_build_get:NNN #1#2#3
33011 { #1 #3 { \if_false: { \fi: \exp_after:wN \__tl_build_get:w #2 } } }
33012 \cs_new:Npn \__tl_build_get:w #1 \__tl_build_last:NNn #2#3#4
33013 {
33014   \exp_not:n {#4}

```

```

33015     \if_meaning:w \c_empty_tl #3
33016     \exp_after:wN \__tl_build_get_end:w
33017     \fi:
33018     \exp_after:wN \__tl_build_get:w #3
33019   }
33020 \cs_new:Npn \__tl_build_get_end:w #1#2#3
33021 { \exp_after:wN \exp_not:n \exp_after:wN { \if_false: } \fi: }

```

(End definition for `\tl_build_get:NN` and others. This function is documented on page 276.)

`\tl_build_end:N` Get the data then clear the *<next tl>* recursively until finding an empty one. It is perhaps wasteful to repeatedly use `\cs_to_sr:N`. The local/global scope is checked by `\tl_set:Nx` or `\tl_gset:Nx`.

```

33022 \cs_new_protected:Npn \tl_build_end:N #1
33023 {
33024   \__tl_build_get:NNN \__kernel_tl_set:Nx #1 #1
33025   \exp_args:Nc \__tl_build_end_loop:NN { \cs_to_str:N #1 ' } \tl_clear:N
33026 }
33027 \cs_new_protected:Npn \tl_build_gend:N #1
33028 {
33029   \__tl_build_get:NNN \__kernel_tl_gset:Nx #1 #1
33030   \exp_args:Nc \__tl_build_end_loop:NN { \cs_to_str:N #1 ' } \tl_gclear:N
33031 }
33032 \cs_new_protected:Npn \__tl_build_end_loop:NN #1#2
33033 {
33034   \if_meaning:w \c_empty_tl #1
33035   \exp_after:wN \use_none:nnnnnn
33036   \fi:
33037   #2 #1
33038   \exp_args:Nc \__tl_build_end_loop:NN { \cs_to_str:N #1 ' } #2
33039 }

```

(End definition for `\tl_build_end:N`, `\tl_build_gend:N`, and `__tl_build_end_loop:NN`. These functions are documented on page 276.)

51.9.2 Other additions to `\l3tl`

`\tl_range_braced:Nnn` For the braced version `__tl_range_braced:w` sets up `__tl_range_collect_braced:w` which stores items one by one in an argument after the semicolon. The unbraced version is almost identical. The version preserving braces and spaces starts by deleting spaces before the argument to avoid collecting them, and sets up `__tl_range_collect:nn` with a first argument of the form `{ {<collected>} <tokens> }`, whose head is the collected tokens and whose tail is what remains of the original token list. This form makes it easier to move tokens to the *<collected>* tokens.

```

\__tl_range_braced:w
\__tl_range_collect_braced:w
\__tl_range_unbraced:w
\__tl_range_collect_unbraced:w
33040 \cs_new:Npn \tl_range_braced:Nnn { \exp_args:No \tl_range_braced:nnn }
33041 \cs_generate_variant:Nn \tl_range_braced:Nnn { c }
33042 \cs_new:Npn \tl_range_braced:nnn { \__tl_range:Nnnn \__tl_range_braced:w }
33043 \cs_new:Npn \tl_range_unbraced:Nnn
33044 { \exp_args:No \tl_range_unbraced:nnn }
33045 \cs_generate_variant:Nn \tl_range_unbraced:Nnn { c }
33046 \cs_new:Npn \tl_range_unbraced:nnn
33047 { \__tl_range:Nnnn \__tl_range_unbraced:w }
33048 \cs_new:Npn \__tl_range_braced:w #1 ; #2
33049 { \__tl_range_collect_braced:w #1 ; { } #2 }

```

```

33050 \cs_new:Npn \__tl_range_unbraced:w #1 ; #2
33051 { \__tl_range_collect_unbraced:w #1 ; { } #2 }
33052 \cs_new:Npn \__tl_range_collect_braced:w #1 ; #2#3
33053 {
33054   \if_int_compare:w #1 > 1 \exp_stop_f:
33055     \exp_after:wN \__tl_range_collect_braced:w
33056     \int_value:w \int_eval:n { #1 - 1 } \exp_after:wN ;
33057   \fi:
33058   { #2 {#3} }
33059 }
33060 \cs_new:Npn \__tl_range_collect_unbraced:w #1 ; #2#3
33061 {
33062   \if_int_compare:w #1 > 1 \exp_stop_f:
33063     \exp_after:wN \__tl_range_collect_unbraced:w
33064     \int_value:w \int_eval:n { #1 - 1 } \exp_after:wN ;
33065   \fi:
33066   { #2 #3 }
33067 }

```

(End definition for `\tl_range_braced:Nnn` and others. These functions are documented on page 275.)

51.10 Additions to `l3token`

`\c_catcode_active_tl` While `\char_generate:nn` can produce active characters in some engines it cannot in general. It would be possible to simply change the catcode of space but then the code would need to avoid all spaces, making it quite unreadable. Instead we use the primitive `\tex_lowercase:D` trick.

```

33068 \group_begin:
33069   \char_set_catcode_active:N *
33070   \char_set_lccode:nn { '*' } { '\ }
33071   \tex_lowercase:D { \tl_const:Nn \c_catcode_active_space_tl { * } }
33072 \group_end:

```

(End definition for `\c_catcode_active_space_tl`. This variable is documented on page 276.)

```

33073 <@@=peek>

```

```

\l__peek_collect_tl

```

```

33074 \tl_new:N \l__peek_collect_tl

```

(End definition for `\l__peek_collect_tl`.)

`\peek_catcode_collect_inline:Nn`
`\peek_charcode_collect_inline:Nn`
`\peek_meaning_collect_inline:Nn`
`__peek_collect:NNn`
`__peek_collect_true:w`
`__peek_collect_remove:nw`
`__peek_collect:N`

Most of the work is done by `__peek_execute_branches_...`, which calls either `__peek_true:w` or `__peek_false:w` according to whether the next token `\l_peek_token` matches the search token (stored in `\l__peek_search_token` and `\l__peek_search_tl`). Here, in the true case we run `__peek_collect_true:w`, which generally calls `__peek_collect:N` to store the peeked token into `\l__peek_collect_tl`, except in special non-N-type cases (begin-group, end-group, or space), where a frozen token is stored. The true branch calls `__peek_execute_branches_...` to fetch more matching tokens. Once there are no more, `__peek_false_aux:n` closes the safe-align group and runs the user's inline code.

```

33075 \cs_new_protected:Npn \peek_catcode_collect_inline:Nn
33076 { \__peek_collect:NNn \__peek_execute_branches_catcode: }
33077 \cs_new_protected:Npn \peek_charcode_collect_inline:Nn

```

```

33078 { \__peek_collect:NNn \__peek_execute_branches_charcode: }
33079 \cs_new_protected:Npn \peek_meaning_collect_inline:Nn
33080 { \__peek_collect:NNn \__peek_execute_branches_meaning: }
33081 \cs_new_protected:Npn \__peek_collect:NNn #1#2#3
33082 {
33083   \group_align_safe_begin:
33084   \cs_set_eq:NN \l__peek_search_token #2
33085   \tl_set:Nn \l__peek_search_tl {#2}
33086   \tl_clear:N \l__peek_collect_tl
33087   \cs_set:Npn \__peek_false:w
33088     { \exp_args:No \__peek_false_aux:n \l__peek_collect_tl }
33089   \cs_set:Npn \__peek_false_aux:n ##1
33090     {
33091       \group_align_safe_end:
33092       #3
33093     }
33094   \cs_set_eq:NN \__peek_true:w \__peek_collect_true:w
33095   \cs_set:Npn \__peek_true_aux:w { \peek_after:Nw #1 }
33096   \__peek_true_aux:w
33097 }
33098 \cs_new_protected:Npn \__peek_collect_true:w
33099 {
33100   \if_case:w
33101     \if_catcode:w \exp_not:N \l_peek_token { 1 \exp_stop_f: \fi:
33102     \if_catcode:w \exp_not:N \l_peek_token } 2 \exp_stop_f: \fi:
33103     \if_meaning:w \l_peek_token \c_space_token 3 \exp_stop_f: \fi:
33104     0 \exp_stop_f:
33105     \exp_after:wN \__peek_collect:N
33106     \or: \__peek_collect_remove:nw { \c_group_begin_token }
33107     \or: \__peek_collect_remove:nw { \c_group_end_token }
33108     \or: \__peek_collect_remove:nw { ~ }
33109     \fi:
33110 }
33111 \cs_new_protected:Npn \__peek_collect:N #1
33112 {
33113   \tl_put_right:Nn \l__peek_collect_tl {#1}
33114   \__peek_true_aux:w
33115 }
33116 \cs_new_protected:Npn \__peek_collect_remove:nw #1
33117 {
33118   \tl_put_right:Nn \l__peek_collect_tl {#1}
33119   \exp_after:wN \__peek_true_remove:w
33120 }

```

(End definition for `\peek_catcode_collect_inline:Nn` and others. These functions are documented on page 277.)

```

33121 \</package>

```

52 l3deprecation implementation

```

33122 \*package>
33123 \*kernel>

```

```
33124 <@@=deprecation>
```

52.1 Helpers and variables

`\l_deprecation_grace_period_bool` This is set to `true` when the deprecated command that is being defined is in its grace period, meaning between the time it becomes an error by default and the time 6 months later where even `undo-recent-deprecations` stops restoring it.

```
33125 \bool_new:N \l_deprecation_grace_period_bool
```

(End definition for `\l_deprecation_grace_period_bool`.)

`\s_deprecation_mark` Internal scan marks.

```
\s_deprecation_stop 33126 \scan_new:N \s_deprecation_mark
```

```
33127 \scan_new:N \s_deprecation_stop
```

(End definition for `\s_deprecation_mark` and `\s_deprecation_stop`.)

`_deprecation_date_compare:nNnTF` Expects `#1` and `#3` to be dates in the format YYYY-MM-DD (but accepts YYYY or YYYY-MM too, filling in zeros for the missing data). Compares them using `#2` (one of `<`, `=`, `>`).

`_deprecation_date_compare_aux:w`

```
33128 \cs_new:Npn \_deprecation_date_compare:nNnTF #1#2#3
```

```
33129 { \_deprecation_date_compare_aux:w #1 -0-0- \s_deprecation_mark #2 #3 -0-0- \s_depreca
```

```
33130 \cs_new:Npn \_deprecation_date_compare_aux:w
```

```
33131 #1 - #2 - #3 - #4 \s_deprecation_mark #5 #6 - #7 - #8 - #9 \s_deprecation_stop
```

```
33132 {
```

```
33133 \int_compare:nNnTF {#1} = {#6}
```

```
33134 {
```

```
33135 \int_compare:nNnTF {#2} = {#7}
```

```
33136 { \int_compare:nNnTF {#3} #5 {#8} }
```

```
33137 { \int_compare:nNnTF {#2} #5 {#7} }
```

```
33138 }
```

```
33139 { \int_compare:nNnTF {#1} #5 {#6} }
```

```
33140 }
```

(End definition for `_deprecation_date_compare:nNnTF` and `_deprecation_date_compare_aux:w`.)

`\g_kernel_deprecation_undo_recent_bool`

```
33141 \bool_new:N \g_kernel_deprecation_undo_recent_bool
```

(End definition for `\g_kernel_deprecation_undo_recent_bool`.)

`_deprecation_not_yet_deprecated:nTF` Receives a deprecation `<date>` and runs the `true` (`false`) branch if the `expl3` date is earlier (later) than `<date>`. If `undo-recent-deprecations` is used we subtract 6 months to the `expl3` date (equivalently add 6 months to the `<date>`). In addition, if the `expl3` date is between `<date>` and `<date>` plus 6 months, `\l_deprecation_grace_period_bool` is set to `true`, otherwise `false`.

`_deprecation_minus_six_months:w`

```
33142 \cs_new_protected:Npn \_deprecation_not_yet_deprecated:nTF #1
```

```
33143 {
```

```
33144 \bool_set_false:N \l_deprecation_grace_period_bool
```

```
33145 \exp_args:No \_deprecation_date_compare:nNnTF { \ExplLoaderFileDate } < {#1}
```

```
33146 { \use_i:nn }
```

```
33147 {
```

```
33148 \exp_args:Nf \_deprecation_date_compare:nNnTF
```

```
33149 {
```

```
33150 \exp_after:wN \_deprecation_minus_six_months:w
```

```

33151         \ExplLoaderFileDate -0-0- \s__deprecation_stop
33152     } < {#1}
33153     {
33154         \bool_set_true:N \l__deprecation_grace_period_bool
33155         \bool_if:NTF \g__kernel_deprecation_undo_recent_bool
33156     }
33157     { \use_i:nn }
33158 }
33159 }
33160 \cs_new:Npn \__deprecation_minus_six_months:w #1 - #2 - #3 - #4 \s__deprecation_stop
33161 {
33162     \int_compare:nNnTF {#2} > 6
33163     { #1 - \int_eval:n { #2 - 6 } - #3 }
33164     { \int_eval:n { #1 - 1 } - \int_eval:n { #2 + 6 } - #3 }
33165 }

```

(End definition for `__deprecation_not_yet_deprecated:nTF` and `__deprecation_minus_six_months:w`.)

52.2 Patching definitions to deprecate

```

\__kernel_patch_deprecation:nnNNpn {<date>} {<replacement>} {<definition>}
<function> <parameters> {<code>}

```

defines the `<function>` to produce a warning and run its `<code>`, or to produce an error and not run any `<code>`, depending on the `expl3` date.

- If the `expl3` date is less than the `<date>` (plus 6 months in case `undo-recent-deprecations` is used) then we define the `<function>` to produce a warning and run its code. The warning is actually suppressed in two cases:
 - if neither `undo-recent-deprecations` nor `enable-debug` are in effect we may be in an end-user’s document so it is suppressed;
 - if the command is expandable then we cannot produce a warning.
- Otherwise, we define the `<function>` to produce an error.

In both cases we additionally make `\debug_on:n {deprecation}` turn the `<function>` into an `\outer` error, and `\debug_off:n {deprecation}` restore whatever the behaviour was without `\debug_on:n {deprecation}`.

In later sections we use the `l3doc` key `deprecated` with a date equal to that `<date>` plus 6 months, so that `l3doc` will complain if we forget to remove the stale `<parameters>` and `{<code>}`.

In the explanations below, `<definition>` `<function>` `<parameters>` `{<code>}` or assignments that only differ in the scope of the `<definition>` will be called “the standard definition”.

```

\__kernel_patch_deprecation:nnNNpn (The parameter text is grabbed using #5#.) The arguments of \__kernel_deprecation_-
\__deprecation_patch_aux:nnNNnn code:nn are run upon \debug_on:n {deprecation} and \debug_off:n {deprecation},
\__deprecation_warn_once:nnNnn respectively. In both scenarios we the <function> may be \outer so we undefine it with
\__deprecation_patch_aux:Nn \tex_let:D before redefining it, with \__kernel_deprecation_error:Nnn or with some
\__deprecation_just_error:nnNN code added shortly.

```


Then check the date (taking into account `undo-recent-deprecations`) to see if the command should be deprecated right away (false branch of `__deprecation_not_yet_deprecated:nTF`), in which case `__deprecation_just_error:nnNN` makes $\langle function \rangle$ into an error (not `\outer`), ignoring its $\langle parameters \rangle$ and $\langle code \rangle$ completely.

Otherwise distinguish cases where we should give a warning from those where we shouldn't: warnings can only happen for protected commands, and we only want them if either `undo-recent-deprecations` or `enable-debug` is in force, not for standard users.

```

33166 \cs_new_protected:Npn \__kernel_patch_deprecation:nnNNn #1#2#3#4#5#
33167 { \__deprecation_patch_aux:nnNNnn {#1} {#2} #3 #4 {#5} }
33168 \cs_new_protected:Npn \__deprecation_patch_aux:nnNNnn #1#2#3#4#5#6
33169 {
33170   \__kernel_deprecation_code:nn
33171   {
33172     \tex_let:D #4 \scan_stop:
33173     \__kernel_deprecation_error:Nnn #4 {#2} {#1}
33174   }
33175   { \tex_let:D #4 \scan_stop: }
33176   \__deprecation_not_yet_deprecated:nTF {#1}
33177   {
33178     \bool_if:nTF
33179     {
33180       \cs_if_eq_p:NN #3 \cs_gset_protected:Npn &&
33181       \__kernel_if_debug:TF
33182       { \c_true_bool } { \g__kernel_deprecation_undo_recent_bool }
33183     }
33184     { \__deprecation_warn_once:nnNNnn {#1} {#2} #4 {#5} {#6} }
33185     { \__deprecation_patch_aux:Nn #3 { #4 #5 {#6} } }
33186   }
33187   { \__deprecation_just_error:nnNN {#1} {#2} #3 #4 }
33188 }

```

In case we want a warning, the $\langle function \rangle$ is defined to produce such a warning without grabbing any argument, then redefine itself to the standard definition that the $\langle function \rangle$ should have, with arguments, and call that definition. The `x-type` expansion and `\exp_not:n` avoid needing to double the #, which we could not do anyways. We then deal with the code for `\debug_off:n {deprecation}`: presumably someone doing that does not need the warning so we simply do the standard definition.

```

33189 \cs_new_protected:Npn \__deprecation_warn_once:nnNNnn #1#2#3#4#5
33190 {
33191   \cs_gset_protected:Npx #3
33192   {
33193     \__kernel_if_debug:TF
33194     {
33195       \exp_not:N \__kernel_msg_warning:nnxxx
33196       { kernel } { deprecated-command }
33197       {#1}
33198       { \token_to_str:N #3 }
33199       { \tl_to_str:n {#2} }
33200     }
33201     { }
33202     \exp_not:n { \cs_gset_protected:Npn #3 #4 {#5} }
33203     \exp_not:N #3
33204   }

```

```

33205     \__kernel_deprecation_code:nn { }
33206     { \cs_set_protected:Npn #3 #4 {#5} }
33207 }

```

In case we want neither warning nor error, the $\langle function \rangle$ is given its standard definition. Here #1 is $\backslash\text{cs_new:Npn}$ or $\backslash\text{cs_new_protected:Npn}$ and #2 is $\langle function \rangle \langle parameters \rangle \{\langle code \rangle\}$, so #1#2 performs the assignment. For $\backslash\text{debug_off:n}\{\text{deprecation}\}$ we want to use the same assignment but with a different scope, hence the $\backslash\text{cs_if_eq:NNTF}$ test.

```

33208 \cs_new_protected:Npn \__deprecation_patch_aux:Nn #1#2
33209 {
33210     #1 #2
33211     \cs_if_eq:NNTF #1 \cs_gset_protected:Npn
33212     { \__kernel_deprecation_code:nn { } } { \cs_set_protected:Npn #2 } }
33213     { \__kernel_deprecation_code:nn { } } { \cs_set:Npn #2 } }
33214 }

```

Finally, if we want an error we reuse the same $\backslash\text{__deprecation_patch_aux:Nn}$ as the previous case. Indeed, we want $\backslash\text{debug_off:n}\{\text{deprecation}\}$ to make the $\langle function \rangle$ into an error, just like it is by default. The error is expandable or not, and the last argument of the error message is empty or is grace to denote the case where we are in the 6 month grace period, in which case the error message is more detailed.

```

33215 \cs_new_protected:Npn \__deprecation_just_error:nnNN #1#2#3#4
33216 {
33217     \exp_args:NNx \__deprecation_patch_aux:Nn #3
33218     {
33219         \exp_not:N #4
33220         {
33221             \cs_if_eq:NNTF #3 \cs_gset_protected:Npn
33222             { \exp_not:N \__kernel_msg_error:nnnnnn }
33223             { \exp_not:N \__kernel_msg_expandable_error:nnnnnn }
33224             { kernel } { deprecated-command }
33225             {#1}
33226             { \token_to_str:N #4 }
33227             { \tl_to_str:n {#2} }
33228             { \bool_if:NT \l__deprecation_grace_period_bool { grace } }
33229         }
33230     }
33231 }

```

(End definition for $\backslash\text{__kernel_patch_deprecation:nnNNpn}$ and others.)

$\backslash\text{__kernel_deprecation_error:Nnn}$

The $\backslash\text{outer}$ definition here ensures the command cannot appear in an argument. Use this auxiliary on all commands that have been removed since 2015.

```

33232 \cs_new_protected:Npn \__kernel_deprecation_error:Nnn #1#2#3
33233 {
33234     \tex_protected:D \tex_outer:D \tex_edef:D #1
33235     {
33236         \exp_not:N \__kernel_msg_expandable_error:nnnnn
33237         { kernel } { deprecated-command }
33238         { \tl_to_str:n {#3} } { \token_to_str:N #1 } { \tl_to_str:n {#2} }
33239         \exp_not:N \__kernel_msg_error:nnxxx
33240         { kernel } { deprecated-command }
33241         { \tl_to_str:n {#3} } { \token_to_str:N #1 } { \tl_to_str:n {#2} }
33242     }
33243 }

```

(End definition for _kernel_deprecation_error:Nnn.)

```

33244 \_kernel_msg_new:nnn { kernel } { deprecated-command }
33245 {
33246   \tl_if_blank:nF {#3} { Use~ \tl_trim_spaces:n {#3} ~not~ }
33247   #2~deprecated-on~#1.
33248   \str_if_eq:nnT {#4} { grace }
33249   {
33250     \c_space_tl
33251     For~6~months~after~that~date~one~can~restore~a~deprecated~
33252     command~by~loading~the~expl3~package~with~the~option~
33253     'undo-recent-deprecations'.
33254   }
33255 }

```

52.3 Removed functions

_deprecation_old_protected:Nnn Short-hands for old commands whose definition does not matter anymore, i.e., commands past the grace period.

```

\_deprecation_old:Nnn
\_deprecation_old:Nnn
33256 \cs_new_protected:Npn \_deprecation_old_protected:Nnn #1#2#3
33257 {
33258   \_kernel_patch_deprecation:nnNNpn {#3} {#2}
33259   \cs_gset_protected:Npn #1 { }
33260 }
33261 \cs_new_protected:Npn \_deprecation_old:Nnn #1#2#3
33262 {
33263   \_kernel_patch_deprecation:nnNNpn {#3} {#2}
33264   \cs_gset:Npn #1 { }
33265 }
33266 \_deprecation_old:Nnn \box_resize:Nnn
33267 { \box_resize_to_wd_and_ht_plus_dp:Nnn } { 2019-01-01 }
33268 \_deprecation_old:Nnn \box_use_clear:N
33269 { \box_use_drop:N } { 2019-01-01 }
33270 \_deprecation_old:Nnn \c_job_name_tl
33271 { \c_sys_jobname_str } { 2017-01-01 }
33272 \_deprecation_old:Nnn \c_minus_one
33273 { -1 } { 2019-01-01 }
33274 \_deprecation_old:Nnn \c_zero
33275 { 0 } { 2020-01-01 }
33276 \_deprecation_old:Nnn \c_one
33277 { 1 } { 2020-01-01 }
33278 \_deprecation_old:Nnn \c_two
33279 { 2 } { 2020-01-01 }
33280 \_deprecation_old:Nnn \c_three
33281 { 3 } { 2020-01-01 }
33282 \_deprecation_old:Nnn \c_four
33283 { 4 } { 2020-01-01 }
33284 \_deprecation_old:Nnn \c_five
33285 { 5 } { 2020-01-01 }
33286 \_deprecation_old:Nnn \c_six
33287 { 6 } { 2020-01-01 }
33288 \_deprecation_old:Nnn \c_seven
33289 { 7 } { 2020-01-01 }
33290 \_deprecation_old:Nnn \c_eight

```

```

33291 { 8 } { 2020-01-01 }
33292 \_deprecation_old:Nnn \c_nine
33293 { 9 } { 2020-01-01 }
33294 \_deprecation_old:Nnn \c_ten
33295 { 10 } { 2020-01-01 }
33296 \_deprecation_old:Nnn \c_eleven
33297 { 11 } { 2020-01-01 }
33298 \_deprecation_old:Nnn \c_twelve
33299 { 12 } { 2020-01-01 }
33300 \_deprecation_old:Nnn \c_thirteen
33301 { 13 } { 2020-01-01 }
33302 \_deprecation_old:Nnn \c_fourteen
33303 { 14 } { 2020-01-01 }
33304 \_deprecation_old:Nnn \c_fifteen
33305 { 15 } { 2020-01-01 }
33306 \_deprecation_old:Nnn \c_sixteen
33307 { 16 } { 2020-01-01 }
33308 \_deprecation_old:Nnn \c_thirty_two
33309 { 32 } { 2020-01-01 }
33310 \_deprecation_old:Nnn \c_one_hundred
33311 { 100 } { 2020-01-01 }
33312 \_deprecation_old:Nnn \c_two_hundred_fifty_five
33313 { 255 } { 2020-01-01 }
33314 \_deprecation_old:Nnn \c_two_hundred_fifty_six
33315 { 256 } { 2020-01-01 }
33316 \_deprecation_old:Nnn \c_one_thousand
33317 { 1000 } { 2020-01-01 }
33318 \_deprecation_old:Nnn \c_ten_thousand
33319 { 10000 } { 2020-01-01 }
33320 \_deprecation_old:Nnn \dim_case:nnn
33321 { \dim_case:nnF } { 2015-07-14 }
33322 \_deprecation_old:Nnn \file_add_path:nN
33323 { \file_get_full_name:nN } { 2019-01-01 }
33324 \_deprecation_old_protected:Nnn \file_if_exist_input:nT
33325 { \file_if_exist:nT and~ \file_input:n } { 2018-03-05 }
33326 \_deprecation_old_protected:Nnn \file_if_exist_input:nTF
33327 { \file_if_exist:nT and~ \file_input:n } { 2018-03-05 }
33328 \_deprecation_old:Nnn \file_list:
33329 { \file_log_list: } { 2019-01-01 }
33330 \_deprecation_old:Nnn \file_path_include:n
33331 { \seq_put_right:Nn \l_file_search_path_seq } { 2019-01-01 }
33332 \_deprecation_old:Nnn \file_path_remove:n
33333 { \seq_remove_all:Nn \l_file_search_path_seq } { 2019-01-01 }
33334 \_deprecation_old:Nnn \g_file_current_name_tl
33335 { \g_file_curr_name_str } { 2019-01-01 }
33336 \_deprecation_old:Nnn \int_case:nnn
33337 { \int_case:nnF } { 2015-07-14 }
33338 \_deprecation_old:Nnn \int_from_binary:n
33339 { \int_from_bin:n } { 2016-01-05 }
33340 \_deprecation_old:Nnn \int_from_hexadecimal:n
33341 { \int_from_hex:n } { 2016-01-05 }
33342 \_deprecation_old:Nnn \int_from_octal:n
33343 { \int_from_oct:n } { 2016-01-05 }
33344 \_deprecation_old:Nnn \int_to_binary:n

```

```

33345 { \int_to_bin:n } { 2016-01-05 }
33346 \__deprecation_old:Nnn \int_to_hexadecimal:n
33347 { \int_to_hex:n } { 2016-01-05 }
33348 \__deprecation_old:Nnn \int_to_octal:n
33349 { \int_to_oct:n } { 2016-01-05 }
33350 \__deprecation_old_protected:Nnn \ior_get_str:NN
33351 { \ior_str_get:NN } { 2018-03-05 }
33352 \__deprecation_old:Nnn \ior_list_streams:
33353 { \ior_show_list: } { 2019-01-01 }
33354 \__deprecation_old:Nnn \ior_log_streams:
33355 { \ior_log_list: } { 2019-01-01 }
33356 \__deprecation_old:Nnn \iow_list_streams:
33357 { \iow_show_list: } { 2019-01-01 }
33358 \__deprecation_old:Nnn \iow_log_streams:
33359 { \iow_log_list: } { 2019-01-01 }
33360 \__deprecation_old:Nnn \lua_escape_x:n
33361 { \lua_escape:e } { 2020-01-01 }
33362 \__deprecation_old:Nnn \lua_now_x:n
33363 { \lua_now:e } { 2020-01-01 }
33364 \__deprecation_old_protected:Nnn \lua_shipout_x:n
33365 { \lua_shipout_e:n } { 2020-01-01 }
33366 \__deprecation_old:Nnn \luatex_if_engine_p:
33367 { \sys_if_engine luatex_p: } { 2017-01-01 }
33368 \__deprecation_old:Nnn \luatex_if_engine:F
33369 { \sys_if_engine luatex:F } { 2017-01-01 }
33370 \__deprecation_old:Nnn \luatex_if_engine:T
33371 { \sys_if_engine luatex:T } { 2017-01-01 }
33372 \__deprecation_old:Nnn \luatex_if_engine:TF
33373 { \sys_if_engine luatex:TF } { 2017-01-01 }
33374 \__deprecation_old_protected:Nnn \msg_interrupt:nnn
33375 { [Defined-error-message] } { 2020-01-01 }
33376 \__deprecation_old_protected:Nnn \msg_log:n
33377 { \iow_log:n } { 2020-01-01 }
33378 \__deprecation_old_protected:Nnn \msg_term:n
33379 { \iow_term:n } { 2020-01-01 }
33380 \__deprecation_old:Nnn \pdftex_if_engine_p:
33381 { \sys_if_engine pdftex_p: } { 2017-01-01 }
33382 \__deprecation_old:Nnn \pdftex_if_engine:F
33383 { \sys_if_engine pdftex:F } { 2017-01-01 }
33384 \__deprecation_old:Nnn \pdftex_if_engine:T
33385 { \sys_if_engine pdftex:T } { 2017-01-01 }
33386 \__deprecation_old:Nnn \pdftex_if_engine:TF
33387 { \sys_if_engine pdftex:TF } { 2017-01-01 }
33388 \__deprecation_old:Nnn \prop_get:cn
33389 { \prop_item:cn } { 2016-01-05 }
33390 \__deprecation_old:Nnn \prop_get:Nn
33391 { \prop_item:Nn } { 2016-01-05 }
33392 \__deprecation_old:Nnn \quark_if_recursion_tail_break:N
33393 { } { 2015-07-14 }
33394 \__deprecation_old:Nnn \quark_if_recursion_tail_break:n
33395 { } { 2015-07-14 }
33396 \__deprecation_old:Nnn \scan_align_safe_stop:
33397 { protected-commands } { 2017-01-01 }
33398 \__deprecation_old:Nnn \sort_ordered:

```

```

33399 { \sort_return_same: } { 2019-01-01 }
33400 \__deprecation_old:Nnn \sort_reversed:
33401 { \sort_return_swapped: } { 2019-01-01 }
33402 \__deprecation_old:Nnn \str_case:nnn
33403 { \str_case:nnF } { 2015-07-14 }
33404 \__deprecation_old:Nnn \str_case:onn
33405 { \str_case:onF } { 2015-07-14 }
33406 \__deprecation_old:Nnn \str_case_x:nn
33407 { \str_case_e:nn } { 2020-01-01 }
33408 \__deprecation_old:Nnn \str_case_x:nnn
33409 { \str_case_e:nnF } { 2015-07-14 }
33410 \__deprecation_old:Nnn \str_case_x:nnT
33411 { \str_case_e:nnT } { 2020-01-01 }
33412 \__deprecation_old:Nnn \str_case_x:nnTF
33413 { \str_case_e:nnTF } { 2020-01-01 }
33414 \__deprecation_old:Nnn \str_case_x:nnF
33415 { \str_case_e:nnF } { 2020-01-01 }
33416 \__deprecation_old:Nnn \str_if_eq_x_p:nn
33417 { \str_if_eq_p:ee } { 2020-01-01 }
33418 \__deprecation_old:Nnn \str_if_eq_x:nnT
33419 { \str_if_eq:eeT } { 2020-01-01 }
33420 \__deprecation_old:Nnn \str_if_eq_x:nnF
33421 { \str_if_eq:eeF } { 2020-01-01 }
33422 \__deprecation_old:Nnn \str_if_eq_x:nnTF
33423 { \str_if_eq:eeTF } { 2020-01-01 }
33424 \__deprecation_old_protected:Nnn \tl_show_analysis:n
33425 { \tl_analysis_show:N } { 2020-01-01 }
33426 \__deprecation_old_protected:Nnn \tl_show_analysis:n
33427 { \tl_analysis_show:n } { 2020-01-01 }
33428 \__deprecation_old:Nnn \tl_case:cnn
33429 { \tl_case:cnF } { 2015-07-14 }
33430 \__deprecation_old:Nnn \tl_case:Nnn
33431 { \tl_case:NnF } { 2015-07-14 }
33432 \__deprecation_old_protected:Nnn \tl_to_lowercase:n
33433 { \tex_lowercase:D } { 2018-03-05 }
33434 \__deprecation_old_protected:Nnn \tl_to_uppercase:n
33435 { \tex_uppercase:D } { 2018-03-05 }
33436 \__deprecation_old:Nnn \token_new:Nn
33437 { \cs_new_eq:NN } { 2019-01-01 }
33438 \__deprecation_old:Nnn \xetex_if_engine_p:
33439 { \sys_if_engine_xetex_p: } { 2017-01-01 }
33440 \__deprecation_old:Nnn \xetex_if_engine:F
33441 { \sys_if_engine_xetex:F } { 2017-01-01 }
33442 \__deprecation_old:Nnn \xetex_if_engine:T
33443 { \sys_if_engine_xetex:T } { 2017-01-01 }
33444 \__deprecation_old:Nnn \xetex_if_engine:TF
33445 { \sys_if_engine_xetex:TF } { 2017-01-01 }

```

(End definition for __deprecation_old_protected:Nnn and __deprecation_old:Nnn.)

52.4 Loading the patches

When loaded first, the patches are simply read here. Here the deprecation code is loaded with the lower-level __kernel_... macro because we don't want it to flip the \g__-

`sys_deprecation_bool` boolean, so that the deprecation code can be re-loaded later (when using `undo-recent-deprecations`).

```

33446 \group_begin:
33447 \cs_set_protected:Npn \ProvidesExplFile
33448 {
33449   \char_set_catcode_space:n { '\ }
33450   \ProvidesExplFileAux
33451 }
33452 \cs_set_protected:Npx \ProvidesExplFileAux #1#2#3#4
33453 {
33454   \group_end:
33455   \cs_if_exist:NTF \ProvidesFile
33456     { \exp_not:N \ProvidesFile {#1} [ #2~v#3~#4 ] }
33457     { \iow_log:x { File:~#1~#2~v#3~#4 } }
33458 }
33459 \cs_gset_protected:Npn \__kernel_sys_configuration_load:n #1
33460 { \file_input:n { #1 .def } }
33461 \__kernel_sys_configuration_load:n { l3deprecation }
33462 </kernel>
33463 <*patches>

Standard file identification.
33464 \ProvidesExplFile{l3deprecation.def}{2019-04-06}{L3 Deprecated functions}

```

52.5 Deprecated l3box functions

```

\box_set_eq_clear:NN
\box_set_eq_clear:cN
\box_set_eq_clear:Nc
\box_set_eq_clear:cc
\box_gset_eq_clear:NN
\box_gset_eq_clear:cN
\box_gset_eq_clear:Nc
\box_gset_eq_clear:cc
33465 \__kernel_patch_deprecation:nnNNpn { 2021-01-01 } { \box_set_eq_drop:N }
33466 \cs_gset_protected:Npn \box_set_eq_clear:NN #1#2
33467 { \tex_setbox:D #1 \tex_box:D #2 }
33468 \__kernel_patch_deprecation:nnNNpn { 2021-01-01 } { \box_gset_eq_drop:N }
33469 \cs_gset_protected:Npn \box_gset_eq_clear:NN #1#2
33470 { \tex_global:D \tex_setbox:D #1 \tex_box:D #2 }
33471 \cs_generate_variant:Nn \box_set_eq_clear:NN { c , Nc , cc }
33472 \cs_generate_variant:Nn \box_gset_eq_clear:NN { c , Nc , cc }

(End definition for \box_set_eq_clear:NN and \box_gset_eq_clear:NN.)

```

```

\hbox_unpack_clear:N
\hbox_unpack_clear:c
33473 \__kernel_patch_deprecation:nnNNpn { 2021-01-01 } { \hbox_unpack_drop:N }
33474 \cs_gset_protected:Npn \hbox_unpack_clear:N
33475 { \hbox_unpack_drop:N }
33476 \cs_generate_variant:Nn \hbox_unpack_clear:N { c }

(End definition for \hbox_unpack_clear:N.)

```

```

\vbox_unpack_clear:N
\vbox_unpack_clear:c
33477 \__kernel_patch_deprecation:nnNNpn { 2021-01-01 } { \vbox_unpack_drop:N }
33478 \cs_gset_protected:Npn \vbox_unpack_clear:N
33479 { \vbox_unpack_drop:N }
33480 \cs_generate_variant:Nn \vbox_unpack_clear:N { c }

(End definition for \vbox_unpack_clear:N.)

```

52.6 Deprecated l3str functions

33481 <@@=str>

```
\str_lower_case:n
\str_lower_case:f
\str_upper_case:n
\str_upper_case:f
\str_fold_case:n
\str_fold_case:V
33482 \__kernel_patch_deprecation:nnNNpn { 2022-01-01 } { \str_lowercase:n }
33483 \cs_gset:Npn \str_lower_case:n { \str_lowercase:n }
33484 \__kernel_patch_deprecation:nnNNpn { 2022-01-01 } { \str_lowercase:f }
33485 \cs_gset:Npn \str_lower_case:f { \str_lowercase:f }
33486 \__kernel_patch_deprecation:nnNNpn { 2022-01-01 } { \str_uppercase:n }
33487 \cs_gset:Npn \str_upper_case:n { \str_uppercase:n }
33488 \__kernel_patch_deprecation:nnNNpn { 2022-01-01 } { \str_uppercase:f }
33489 \cs_gset:Npn \str_upper_case:f { \str_uppercase:f }
33490 \__kernel_patch_deprecation:nnNNpn { 2022-01-01 } { \str_foldcase:n }
33491 \cs_gset:Npn \str_fold_case:n { \str_foldcase:n }
33492 \__kernel_patch_deprecation:nnNNpn { 2022-01-01 } { \str_foldcase:V }
33493 \cs_gset:Npn \str_fold_case:V { \str_foldcase:V }
```

(End definition for \str_lower_case:n, \str_upper_case:n, and \str_fold_case:n.)

\str_declare_eight_bit_encoding:nnn

This command was made internal, with one more argument. There is no easy way to compute a reasonable value for that extra argument so we take a value that is big enough to accomodate all of Unicode.

```
33494 \__kernel_patch_deprecation:nnNNpn { 2022-01-01 } { }
33495 \cs_gset_protected:Npn \str_declare_eight_bit_encoding:nnn #1
33496 { \__str_declare_eight_bit_encoding:nnnn {#1} { 1114112 } }
```

(End definition for \str_declare_eight_bit_encoding:nnn.)

52.7 Deprecated l3seq functions

\seq_indexed_map_inline:Nn

\seq_indexed_map_function:NN

```
33497 \__kernel_patch_deprecation:nnNNpn { 2022-07-01 } { \seq_map_indexed_inline:Nn }
33498 \cs_gset:Npn \seq_indexed_map_inline:Nn { \seq_map_indexed_inline:Nn }
33499 \__kernel_patch_deprecation:nnNNpn { 2022-07-01 } { \seq_map_indexed_function:NN }
33500 \cs_gset:Npn \seq_indexed_map_function:NN { \seq_map_indexed_function:NN }
```

(End definition for \seq_indexed_map_inline:Nn and \seq_indexed_map_function:NN.)

52.7.1 Deprecated l3tl functions

33501 <@@=tl>

```
\tl_set_from_file:Nnn
\tl_set_from_file:cnn
\tl_gset_from_file:Nnn
\tl_gset_from_file:cnn
\tl_set_from_file_x:Nnn
\tl_set_from_file_x:cnn
\tl_gset_from_file_x:Nnn
\tl_gset_from_file_x:cnn
33502 \__kernel_patch_deprecation:nnNNpn { 2021-01-01 } { \file_get:nnN }
33503 \cs_gset_protected:Npn \tl_set_from_file:Nnn #1#2#3
33504 { \file_get:nnN {#3} {#2} #1 }
33505 \cs_generate_variant:Nn \tl_set_from_file:Nnn { c }
33506 \__kernel_patch_deprecation:nnNNpn { 2021-01-01 } { \file_get:nnN }
33507 \cs_gset_protected:Npn \tl_gset_from_file:Nnn #1#2#3
33508 {
33509 \group_begin:
33510 \file_get:nnN {#3} {#2} \l__tl_internal_a_tl
33511 \tl_gset_eq:NN #1 \l__tl_internal_a_tl
33512 \group_end:
```



```

33513 }
33514 \cs_generate_variant:Nn \tl_gset_from_file:Nnn { c }
33515 \__kernel_patch_deprecation:nnNNpn { 2021-01-01 } { \file_get:nnN }
33516 \cs_gset_protected:Npn \tl_set_from_file_x:Nnn #1#2#3
33517 {
33518   \group_begin:
33519     \file_get:nnN {#3} {#2} \l__tl_internal_a_tl
33520     #2 \scan_stop:
33521     \__kernel_tl_set:Nx \l__tl_internal_a_tl { \l__tl_internal_a_tl }
33522     \exp_args:NNNo \group_end:
33523     \tl_set:Nn #1 \l__tl_internal_a_tl
33524   }
33525 \cs_generate_variant:Nn \tl_set_from_file_x:Nnn { c }
33526 \__kernel_patch_deprecation:nnNNpn { 2021-01-01 } { \file_get:nnN }
33527 \cs_gset_protected:Npn \tl_gset_from_file_x:Nnn #1#2#3
33528 {
33529   \group_begin:
33530     \file_get:nnN {#3} {#2} \l__tl_internal_a_tl
33531     #2 \scan_stop:
33532     \__kernel_tl_gset:Nx #1 { \l__tl_internal_a_tl }
33533   \group_end:
33534 }
33535 \cs_generate_variant:Nn \tl_gset_from_file_x:Nnn { c }

```

(End definition for \tl_set_from_file:Nnn and others.)

```

\tl_lower_case:n
\tl_lower_case:nn 33536 \__kernel_patch_deprecation:nnNNpn { 2022-01-01 } { \text_lowercase:n }
\tl_upper_case:n 33537 \cs_gset:Npn \tl_lower_case:n #1
\tl_upper_case:nn 33538 { \text_lowercase:n {#1} }
\tl_mixed_case:n 33539 \__kernel_patch_deprecation:nnNNpn { 2022-01-01 } { \text_lowercase:nn }
\tl_mixed_case:nn 33540 \cs_gset:Npn \tl_lower_case:nn #1#2
33541 { \text_lowercase:nn {#1} {#2} }
33542 \__kernel_patch_deprecation:nnNNpn { 2022-01-01 } { \text_uppercase:n }
33543 \cs_gset:Npn \tl_upper_case:n #1
33544 { \text_uppercase:n {#1} }
33545 \__kernel_patch_deprecation:nnNNpn { 2022-01-01 } { \text_uppercase:nn }
33546 \cs_gset:Npn \tl_upper_case:nn #1#2
33547 { \text_uppercase:nn {#1} {#2} }
33548 \__kernel_patch_deprecation:nnNNpn { 2022-01-01 } { \text_titlecase:n }
33549 \cs_gset:Npn \tl_mixed_case:n #1
33550 { \text_titlecase:n {#1} }
33551 \__kernel_patch_deprecation:nnNNpn { 2022-01-01 } { \text_titlecase:nn }
33552 \cs_gset:Npn \tl_mixed_case:nn #1#2
33553 { \text_titlecase:nn {#1} {#2} }

```

(End definition for \tl_lower_case:n and others.)

52.8 Deprecated l3token functions

```

\token_get_prefix_spec:N
\token_get_arg_spec:N 33554 \__kernel_patch_deprecation:nnNNpn { 2021-01-01 } { \cs_prefix_spec:N }
\token_get_replacement_spec:N 33555 \cs_gset:Npn \token_get_prefix_spec:N { \cs_prefix_spec:N }
33556 \__kernel_patch_deprecation:nnNNpn { 2021-01-01 } { \cs_argument_spec:N }

```

```

33557 \cs_gset:Npn \token_get_arg_spec:N { \cs_argument_spec:N }
33558 \__kernel_patch_deprecation:nnNNpn { 2021-01-01 } { \cs_replacement_spec:N }
33559 \cs_gset:Npn \token_get_replacement_spec:N { \cs_replacement_spec:N }

```

(End definition for \token_get_prefix_spec:N, \token_get_arg_spec:N, and \token_get_replacement_spec:N.)

```

\char_lower_case:N
\char_upper_case:N
\char_mixed_case:Nn
\char_fold_case:N
\char_str_lower_case:N
\char_str_upper_case:N
\char_str_mixed_case:Nn
\char_str_fold_case:N
33560 \__kernel_patch_deprecation:nnNNpn { 2022-01-01 } { \char_lowercase:N }
33561 \cs_gset:Npn \char_lower_case:N { \char_lowercase:N }
33562 \__kernel_patch_deprecation:nnNNpn { 2022-01-01 } { \char_uppercase:N }
33563 \cs_gset:Npn \char_upper_case:N { \char_uppercase:N }
33564 \__kernel_patch_deprecation:nnNNpn { 2022-01-01 } { \char_titlecase:N }
33565 \cs_gset:Npn \char_mixed_case:N { \char_titlecase:N }
33566 \__kernel_patch_deprecation:nnNNpn { 2022-01-01 } { \char_foldcase:N }
33567 \cs_gset:Npn \char_fold_case:N { \char_foldcase:N }
33568 \__kernel_patch_deprecation:nnNNpn { 2022-01-01 } { \char_str_lowercase:N }
33569 \cs_gset:Npn \char_str_lower_case:N { \char_str_lowercase:N }
33570 \__kernel_patch_deprecation:nnNNpn { 2022-01-01 } { \char_str_uppercase:N }
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33573 \cs_gset:Npn \char_str_mixed_case:N { \char_str_titlecase:N }
33574 \__kernel_patch_deprecation:nnNNpn { 2022-01-01 } { \char_str_foldcase:N }
33575 \cs_gset:Npn \char_str_fold_case:N { \char_str_foldcase:N }

```

(End definition for \char_lower_case:N and others.)

52.9 Deprecated l3file functions

```

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```

(End definition for \c_term_ior.)

```

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